# Table of Contents

I. Tutorial

1. Getting Started .............................................................................................................. 1
   1.1. Installation .............................................................................................................. 1
   1.2. Architectural Fundamentals .................................................................................. 1
   1.3. Creating a Database ............................................................................................... 2
   1.4. Accessing a Database ........................................................................................... 3

2. The SQL Language ........................................................................................................ 5
   2.1. Introduction .............................................................................................................. 5
   2.2. Concepts ................................................................................................................ 5
   2.3. Creating a New Table ............................................................................................. 5
   2.4. Populating a Table With Rows ............................................................................. 6
   2.5. Querying a Table .................................................................................................... 7
   2.6. Joins Between Tables ............................................................................................ 9
   2.7. Aggregate Functions ............................................................................................. 11
   2.8. Updates .................................................................................................................. 12
   2.9. Deletions ............................................................................................................... 13

3. Advanced Features ....................................................................................................... 14
   3.1. Introduction ............................................................................................................. 14
   3.2. Views ...................................................................................................................... 14
   3.3. Foreign Keys ........................................................................................................... 14
   3.4. Transactions .......................................................................................................... 15
   3.5. Window Functions ................................................................................................. 17
   3.6. Inheritance ............................................................................................................ 20
   3.7. Conclusion ............................................................................................................. 21

II. The SQL Language ...................................................................................................... 22

4. SQL Syntax .................................................................................................................. 24
   4.1. Lexical Structure .................................................................................................... 24
   4.1.1. Identifiers and Key Words .................................................................................. 24
   4.1.2. Constants .......................................................................................................... 26
       4.1.2.1. String Constants ......................................................................................... 26
       4.1.2.2. String Constants with C-style Escapes ....................................................... 26
       4.1.2.3. String Constants with Unicode Escapes ..................................................... 28
       4.1.2.4. Dollar-quoted String Constants ................................................................ 28
       4.1.2.5. Bit-string Constants .................................................................................... 29
       4.1.2.6. Numeric Constants ..................................................................................... 29
       4.1.2.7. Constants of Other Types ......................................................................... 30
<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>6. Data Manipulation</td>
<td>89</td>
</tr>
<tr>
<td>6.1. Inserting Data</td>
<td>89</td>
</tr>
<tr>
<td>6.2. Updating Data</td>
<td>90</td>
</tr>
<tr>
<td>6.3. Deleting Data</td>
<td>91</td>
</tr>
<tr>
<td>6.4. Returning Data From Modified Rows</td>
<td>91</td>
</tr>
<tr>
<td>7. Queries</td>
<td>93</td>
</tr>
<tr>
<td>7.1. Overview</td>
<td>93</td>
</tr>
<tr>
<td>7.2. Table Expressions</td>
<td>93</td>
</tr>
<tr>
<td>7.2.1. The FROM Clause</td>
<td>94</td>
</tr>
<tr>
<td>7.2.1.1. Joined Tables</td>
<td>94</td>
</tr>
<tr>
<td>7.2.1.2. Table and Column Aliases</td>
<td>98</td>
</tr>
<tr>
<td>7.2.1.3. Subqueries</td>
<td>99</td>
</tr>
<tr>
<td>7.2.1.4. Table Functions</td>
<td>99</td>
</tr>
<tr>
<td>7.2.1.5. LATERAL Subqueries</td>
<td>100</td>
</tr>
<tr>
<td>7.2.2. The WHERE Clause</td>
<td>102</td>
</tr>
<tr>
<td>7.2.3. The GROUP BY and HAVING Clauses</td>
<td>102</td>
</tr>
<tr>
<td>7.2.4. GROUPING SETS, CUBE, and ROLLUP</td>
<td>105</td>
</tr>
<tr>
<td>7.2.5. Window Function Processing</td>
<td>107</td>
</tr>
<tr>
<td>7.3. Select Lists</td>
<td>107</td>
</tr>
<tr>
<td>7.3.1. Select-List Items</td>
<td>108</td>
</tr>
<tr>
<td>7.3.2. Column Labels</td>
<td>108</td>
</tr>
<tr>
<td>7.3.3. DISTINCT</td>
<td>109</td>
</tr>
<tr>
<td>7.4. Combining Queries</td>
<td>109</td>
</tr>
<tr>
<td>7.5. Sorting Rows</td>
<td>110</td>
</tr>
<tr>
<td>7.6. LIMIT and OFFSET</td>
<td>111</td>
</tr>
<tr>
<td>7.7. VALUES Lists</td>
<td>111</td>
</tr>
<tr>
<td>7.8. WITH Queries (Common Table Expressions)</td>
<td>112</td>
</tr>
<tr>
<td>7.8.1. SELECT in WITH</td>
<td>112</td>
</tr>
<tr>
<td>7.8.2. Data-Modifying Statements in WITH</td>
<td>116</td>
</tr>
<tr>
<td>8. Data Types</td>
<td>118</td>
</tr>
<tr>
<td>8.1. Numeric Types</td>
<td>119</td>
</tr>
<tr>
<td>8.1.1. Integer Types</td>
<td>120</td>
</tr>
<tr>
<td>8.1.2. Arbitrary Precision Numbers</td>
<td>120</td>
</tr>
<tr>
<td>8.1.3. Floating-Point Types</td>
<td>122</td>
</tr>
<tr>
<td>8.1.4. Serial Types</td>
<td>123</td>
</tr>
<tr>
<td>8.2. Monetary Types</td>
<td>124</td>
</tr>
<tr>
<td>8.3. Character Types</td>
<td>125</td>
</tr>
<tr>
<td>8.4. Binary Data Types</td>
<td>127</td>
</tr>
<tr>
<td>8.4.1. bytea Hex Format</td>
<td>127</td>
</tr>
</tbody>
</table>
8.16.1. Declaration of Composite Types.................................................................169
8.16.2. Constructing Composite Values.................................................................170
8.16.3. Accessing Composite Types ......................................................................171
8.16.4. Modifying Composite Types .....................................................................172
8.16.5. Using Composite Types in Queries............................................................172
8.16.6. Composite Type Input and Output Syntax ..................................................174

8.17. Range Types .................................................................................................175
  8.17.1. Built-in Range Types ................................................................................175
  8.17.2. Examples ..................................................................................................176
  8.17.3. Inclusive and Exclusive Bounds ...............................................................176
  8.17.4. Infinite (Unbounded) Ranges ..................................................................176
  8.17.5. Range Input/Output ................................................................................177
  8.17.6. Constructing Ranges ...............................................................................178
  8.17.7. Discrete Range Types ............................................................................178
  8.17.8. Defining New Range Types ....................................................................179
  8.17.9. Indexing ..................................................................................................180
  8.17.10. Constraints on Ranges .........................................................................180

8.18. Object Identifier Types ................................................................................181
8.19. pg_lsn Type ..................................................................................................183
8.20. Pseudo-TYPES..............................................................................................183

9. Functions and Operators ..................................................................................185
  9.1. Logical Operators ........................................................................................185
  9.2. Comparison Functions and Operators .........................................................185
  9.3. Mathematical Functions and Operators .......................................................188
  9.4. String Functions and Operators ..................................................................192
    9.4.1. format....................................................................................................207
  9.5. Binary String Functions and Operators .......................................................209
  9.6. Bit String Functions and Operators .............................................................211
  9.7. Pattern Matching ........................................................................................212
    9.7.1. LIKE....................................................................................................212
    9.7.2. SIMILAR TO Regular Expressions ......................................................213
    9.7.3. POSIX Regular Expressions ..................................................................214
      9.7.3.1. Regular Expression Details .............................................................218
      9.7.3.2. Bracket Expressions .....................................................................220
      9.7.3.3. Regular Expression Escapes ..........................................................221
      9.7.3.4. Regular Expression Metasyntax ....................................................223
      9.7.3.5. Regular Expression Matching Rules ............................................225
      9.7.3.6. Limits and Compatibility ..............................................................226
      9.7.3.7. Basic Regular Expressions ............................................................227
  9.8. Data Type Formatting Functions ..................................................................227
  9.9. Date/Time Functions and Operators ............................................................234
    9.9.1. EXTRACT, date_part ............................................................240
    9.9.2. date_trunc..........................................................................................244
    9.9.3. AT TIME ZONE .................................................................244
    9.9.4. Current Date/Time ............................................................................245
    9.9.5. Delaying Execution ............................................................................247
  9.10. Enum Support Functions ............................................................................247
  9.11. Geometric Functions and Operators ..........................................................248
  9.12. Network Address Functions and Operators ................................................253
  9.13. Text Search Functions and Operators ........................................................255
  9.14. XML Functions ........................................................................................260
<table>
<thead>
<tr>
<th>Section</th>
<th>Page</th>
</tr>
</thead>
<tbody>
<tr>
<td>9.14.1.1. xmlcomment</td>
<td>260</td>
</tr>
<tr>
<td>9.14.1.2. xmlconcat</td>
<td>261</td>
</tr>
<tr>
<td>9.14.1.3. xmlelement</td>
<td>261</td>
</tr>
<tr>
<td>9.14.1.4. xmlforest</td>
<td>263</td>
</tr>
<tr>
<td>9.14.1.5. xmlpi</td>
<td>263</td>
</tr>
<tr>
<td>9.14.1.6. xmlroot</td>
<td>263</td>
</tr>
<tr>
<td>9.14.1.7. xmlagg</td>
<td>264</td>
</tr>
<tr>
<td>9.14.2. XML Predicates</td>
<td>265</td>
</tr>
<tr>
<td>9.14.2.1. IS DOCUMENT</td>
<td>265</td>
</tr>
<tr>
<td>9.14.2.2. IS NOT DOCUMENT</td>
<td>265</td>
</tr>
<tr>
<td>9.14.2.3. XML EXISTS</td>
<td>265</td>
</tr>
<tr>
<td>9.14.2.4. xml_is_well_formed</td>
<td>265</td>
</tr>
<tr>
<td>9.14.3. Processing XML</td>
<td>266</td>
</tr>
<tr>
<td>9.14.4. Mapping Tables to XML</td>
<td>267</td>
</tr>
<tr>
<td>9.15. JSON Functions and Operators</td>
<td>271</td>
</tr>
<tr>
<td>9.16. Sequence Manipulation Functions</td>
<td>280</td>
</tr>
<tr>
<td>9.17. Conditional Expressions</td>
<td>282</td>
</tr>
<tr>
<td>9.17.1. CASE</td>
<td>282</td>
</tr>
<tr>
<td>9.17.2. COALESCE</td>
<td>284</td>
</tr>
<tr>
<td>9.17.3. NULLIF</td>
<td>284</td>
</tr>
<tr>
<td>9.17.4. GREATEST and LEAST</td>
<td>284</td>
</tr>
<tr>
<td>9.18. Array Functions and Operators</td>
<td>284</td>
</tr>
<tr>
<td>9.19. Range Functions and Operators</td>
<td>288</td>
</tr>
<tr>
<td>9.20. Aggregate Functions</td>
<td>290</td>
</tr>
<tr>
<td>9.21. Window Functions</td>
<td>299</td>
</tr>
<tr>
<td>9.22. Subquery Expressions</td>
<td>301</td>
</tr>
<tr>
<td>9.22.1. EXISTS</td>
<td>301</td>
</tr>
<tr>
<td>9.22.2. IN</td>
<td>301</td>
</tr>
<tr>
<td>9.22.3. NOT IN</td>
<td>302</td>
</tr>
<tr>
<td>9.22.4. ANY/SOME</td>
<td>302</td>
</tr>
<tr>
<td>9.22.5. ALL</td>
<td>303</td>
</tr>
<tr>
<td>9.22.6. Single-row Comparison</td>
<td>303</td>
</tr>
<tr>
<td>9.23. Row and Array Comparisons</td>
<td>304</td>
</tr>
<tr>
<td>9.23.1. IN</td>
<td>304</td>
</tr>
<tr>
<td>9.23.2. NOT IN</td>
<td>304</td>
</tr>
<tr>
<td>9.23.3. ANY/SOME (array)</td>
<td>305</td>
</tr>
<tr>
<td>9.23.4. ALL (array)</td>
<td>305</td>
</tr>
<tr>
<td>9.23.5. Row Constructor Comparison</td>
<td>305</td>
</tr>
<tr>
<td>9.23.6. Composite Type Comparison</td>
<td>306</td>
</tr>
<tr>
<td>9.24. Set Returning Functions</td>
<td>307</td>
</tr>
<tr>
<td>9.25. System Information Functions</td>
<td>310</td>
</tr>
<tr>
<td>9.26. System Administration Functions</td>
<td>326</td>
</tr>
<tr>
<td>9.26.2. Server Signaling Functions</td>
<td>327</td>
</tr>
<tr>
<td>9.26.3. Backup Control Functions</td>
<td>328</td>
</tr>
<tr>
<td>9.26.4. Recovery Control Functions</td>
<td>330</td>
</tr>
<tr>
<td>9.26.5. Snapshot Synchronization Functions</td>
<td>332</td>
</tr>
<tr>
<td>9.26.6. Replication Functions</td>
<td>333</td>
</tr>
<tr>
<td>9.26.7. Database Object Management Functions</td>
<td>337</td>
</tr>
<tr>
<td>9.26.10. Advisory Lock Functions</td>
<td>341</td>
</tr>
</tbody>
</table>
12. Full Text Search ................................................................................................................ 375
  12.1. Introduction .......................................................................................................... 375
  12.1.1. What Is a Document?................................................................................ 376
  12.1.2. Basic Text Matching ................................................................................ 376
  12.1.3. Configurations ....................................................................................... 378
  12.2. Tables and Indexes............................................................................................ 379
    12.2.1. Searching a Table .............................................................................. 379
    12.2.2. Creating Indexes .............................................................................. 380
  12.3. Controlling Text Search....................................................................................... 381
    12.3.1. Parsing Documents ........................................................................... 381
    12.3.2. Parsing Queries ................................................................................. 382
    12.3.3. Ranking Search Results ..................................................................... 384
    12.3.4. Highlighting Results ........................................................................... 386
  12.4. Additional Features .............................................................................................. 387
    12.4.1. Manipulating Documents .................................................................... 387
    12.4.2. Manipulating Queries ......................................................................... 388
      12.4.2.1. Query Rewriting ....................................................................... 389
    12.4.3. Triggers for Automatic Updates .......................................................... 391
    12.4.4. Gathering Document Statistics .............................................................. 392
  12.5. Parsers .................................................................................................................. 393
  12.6. Dictionaries........................................................................................................... 394
    12.6.1. Stop Words ......................................................................................... 395
    12.6.2. Simple Dictionary ............................................................................... 396
    12.6.3. Synonym Dictionary ............................................................................. 397
    12.6.4. Thesaurus Dictionary .......................................................................... 399
      12.6.4.1. Thesaurus Configuration .............................................................. 400
    12.6.4.2. Thesaurus Example ...................................................................... 400
    12.6.5. Ispell Dictionary ................................................................................... 401
III. Server Administration ................................................................. 454

16. Installation from Source Code ..................................................... 456
   16.1. Short Version ........................................................................ 456
   16.2. Requirements ...................................................................... 456
   16.3. Getting The Source ............................................................. 458
   16.4. Installation Procedure ......................................................... 458
   16.5. Post-Installation Setup ......................................................... 468
   16.5.1. Shared Libraries .............................................................. 468
   16.5.2. Environment Variables ................................................... 469
   16.6. Supported Platforms .......................................................... 470
   16.7. Platform-specific Notes ...................................................... 470
      16.7.1. AIX ............................................................................. 470
         16.7.1.1. GCC Issues .............................................................. 471
         16.7.1.2. Unix-Domain Sockets Broken ................................. 471
         16.7.1.3. Internet Address Issues ......................................... 471
         16.7.1.4. Memory Management ............................................. 472
      References and Resources ....................................................... 473
   16.7.2. Cygwin ........................................................................... 474
   16.7.3. HP-UX .......................................................................... 474
   16.7.4. macOS .......................................................................... 475
   16.7.5. MinGW/Native Windows ............................................... 476
      16.7.5.1. Collecting Crash Dumps on Windows ....................... 476
   16.7.6. SCO OpenServer and SCO UnixWare .............................. 476
      16.7.6.1. Skunkware ............................................................... 476
      16.7.6.2. GNU Make ............................................................. 477
      16.7.6.3. Readline ................................................................. 477
      16.7.6.4. Using the UDK on OpenServer ................................. 477
      16.7.6.5. Reading the PostgreSQL Man Pages ....................... 477
      16.7.6.6. C99 Issues with the 7.1.1b Feature Supplement .......... 478
      16.7.6.7. Threading on UnixWare ........................................... 478
   16.7.7. Solaris ........................................................................... 478
      16.7.7.1. Required Tools ........................................................ 478
      16.7.7.2. Problems with OpenSSL .......................................... 478
      16.7.7.3. configure Complains About a Failed Test Program .... 478
      16.7.7.4. 64-bit Build Sometimes Crashes .............................. 479
      16.7.7.5. Compiling for Optimal Performance ......................... 479
      16.7.7.6. Using DTrace for Tracing PostgreSQL ...................... 479

17. Installation from Source Code on Windows ................................. 481
   17.1. Building with Visual C++ or the Microsoft Windows SDK .... 481
      17.1.1. Requirements ............................................................. 482
      17.1.2. Special Considerations for 64-bit Windows .................. 484
      17.1.3. Building ................................................................. 484
      17.1.4. Cleaning and Installing ............................................. 484
      17.1.5. Running the Regression Tests ..................................... 485
      17.1.6. Building the Documentation ...................................... 485
   17.2. Building libpq with Visual C++ or Borland C++ ................. 486
      17.2.1. Generated Files ........................................................ 486

18. Server Setup and Operation ...................................................... 488
18.1. The PostgreSQL User Account ................................................................. 488
18.2. Creating a Database Cluster ................................................................. 488
  18.2.1. Use of Secondary File Systems ...................................................... 489
  18.2.2. Use of Network File Systems ....................................................... 489
18.3. Starting the Database Server ............................................................... 490
  18.3.1. Server Start-up Failures ............................................................ 492
  18.3.2. Client Connection Problems ....................................................... 492
18.4. Managing Kernel Resources .............................................................. 493
  18.4.1. Shared Memory and Semaphores ............................................... 493
  18.4.2. systemd RemoveIPC .............................................................. 498
  18.4.3. Resource Limits ......................................................................... 499
  18.4.4. Linux Memory Overcommit ....................................................... 500
  18.4.5. Linux Huge Pages ...................................................................... 501
18.5. Shutting Down the Server .................................................................. 502
18.6. Upgrading a PostgreSQL Cluster ...................................................... 503
  18.6.1. Upgrading Data via pg_dumpall .................................................. 504
  18.6.2. Upgrading Data via pg_upgrade ................................................ 505
  18.6.3. Upgrading Data via Replication ................................................ 505
18.7. Preventing Server Spoofing ............................................................... 506
18.8. Encryption Options ........................................................................... 506
18.9. Secure TCP/IP Connections with SSL ............................................. 507
  18.9.1. Using Client Certificates ........................................................... 508
  18.9.2. SSL Server File Usage ........................................................... 509
  18.9.3. Creating Certificates ............................................................... 509
18.10. Secure TCP/IP Connections with SSH Tunnels ......................... 511
18.11. Registering Event Log on Windows .............................................. 512
19. Server Configuration ........................................................................... 513
  19.1. Setting Parameters ......................................................................... 513
    19.1.1. Parameter Names and Values .................................................. 513
    19.1.2. Parameter Interaction via the Configuration File ..................... 513
    19.1.3. Parameter Interaction via SQL ............................................... 514
    19.1.4. Parameter Interaction via the Shell ........................................ 515
    19.1.5. Managing Configuration File Contents .................................. 515
  19.2. File Locations ................................................................................ 517
  19.3. Connections and Authentication .................................................... 517
    19.3.1. Connection Settings .............................................................. 518
    19.3.2. Security and Authentication ................................................... 520
  19.4. Resource Consumption ............................................................... 522
    19.4.1. Memory ............................................................................. 522
    19.4.2. Disk .................................................................................. 525
    19.4.3. Kernel Resource Usage ......................................................... 525
    19.4.4. Cost-based Vacuum Delay .................................................... 525
    19.4.5. Background Writer .............................................................. 526
    19.4.6. Asynchronous Behavior ......................................................... 527
  19.5. Write Ahead Log ........................................................................... 529
    19.5.1. Settings ............................................................................. 529
    19.5.2. Checkpoints ....................................................................... 533
    19.5.3. Archiving ........................................................................... 534
  19.6. Replication .................................................................................... 535
    19.6.1. Sending Server(s) ................................................................. 535
    19.6.2. Master Server .................................................................... 536
    19.6.3. Standby Servers ................................................................. 537
V. Server Programming ................................................................................................................. 982

36. Extending SQL .................................................................................................................. 984

36.1. How Extensibility Works ............................................................................................. 984

36.2. The PostgreSQL Type System .................................................................................... 984

36.2.1. Base Types .................................................................................................................. 984

36.2.2. Composite Types ....................................................................................................... 984

36.2.3. Domains ..................................................................................................................... 985

36.2.4. Pseudo-Types ............................................................................................................ 985

36.2.5. Polymorphic Types ................................................................................................... 985

36.3. User-defined Functions ............................................................................................... 986

36.4. Query Language (SQL) Functions ................................................................................. 986

36.4.1. Arguments for SQL Functions .................................................................. 987

36.4.2. SQL Functions on Base Types .......................................................................... 988

36.4.3. SQL Functions on Composite Types ...................................................................... 990

36.4.4. SQL Functions with Output Parameters ................................................................. 992

36.4.5. SQL Functions with Variable Numbers of Arguments ................................................. 993

36.4.6. SQL Functions with Default Values for Arguments ................................................... 994
41.6. Control Structures................................................................................................................ 1133
  41.6.1. Returning From a Function ............................................................................................ 1133
    41.6.1.1. RETURN ........................................................................................................... 1133
    41.6.1.2. RETURN NEXT and RETURN QUERY ...................................................... 1133
  41.6.2. Conditionals .................................................................................................................... 1135
    41.6.2.1. IF-THEN.......................................................................................................... 1135
    41.6.2.2. IF-THEN-ELSE .......................................................................................... 1136
    41.6.2.3. IF-THEN-ELSIF ....................................................................................... 1136
    41.6.2.4. Simple CASE ............................................................................................. 1137
    41.6.2.5. Searched CASE ......................................................................................... 1138
  41.6.3. Simple Loops ................................................................................................................... 1138
    41.6.3.1. LOOP ........................................................................................................... 1138
    41.6.3.2. EXIT ........................................................................................................... 1139
    41.6.3.3. CONTINUE ................................................................................................. 1139
    41.6.3.4. WHILE ........................................................................................................ 1140
    41.6.3.5. FOR (Integer Variant) ..................................................................................... 1140
  41.6.4. Looping Through Query Results .................................................................................. 1141
  41.6.5. Looping Through Arrays ............................................................................................ 1142
  41.6.6. Trapping Errors ............................................................................................................ 1143
    41.6.6.1. Obtaining Information About an Error ......................................................... 1145
    41.6.6.2. Obtaining Execution Location Information .................................................. 1147
  41.7. Cursors.......................................................................................................................... 1147
    41.7.1. Declaring Cursor Variables ....................................................................................... 1147
    41.7.2. Opening Cursors ..................................................................................................... 1148
      41.7.2.1. OPEN FOR query ................................................................................... 1148
      41.7.2.2. OPEN FOR EXECUTE ........................................................................... 1149
      41.7.2.3. Opening a Bound Cursor .......................................................................... 1149
    41.7.3. Using Cursors ......................................................................................................... 1150
      41.7.3.1. FETCH ...................................................................................................... 1150
      41.7.3.2. MOVE ...................................................................................................... 1150
      41.7.3.3. UPDATE/DELETE WHERE CURRENT OF ............................................. 1151
      41.7.3.4. CLOSE ..................................................................................................... 1151
      41.7.3.5. Returning Cursors ............................................................................... 1151
    41.7.4. Looping Through a Cursor’s Result .......................................................................... 1153
  41.8. Errors and Messages ....................................................................................................... 1153
    41.8.1. Reporting Errors and Messages ............................................................................... 1153
    41.8.2. Checking Assertions ............................................................................................... 1155
  41.9. Trigger Procedures ......................................................................................................... 1155
    41.9.1. Triggers on Data Changes ...................................................................................... 1156
    41.9.2. Triggers on Events ............................................................................................... 1162
  41.10. PL/pgSQL Under the Hood ............................................................................................ 1163
    41.10.1. Variable Substitution ............................................................................................ 1163
    41.10.2. Plan Caching ....................................................................................................... 1165
  41.11. Tips for Developing in PL/pgSQL .................................................................................... 1167
    41.11.1. Handling of Quotation Marks .............................................................................. 1167
    41.11.2. Additional Compile-time Checks ........................................................................ 1169
  41.12. Porting from Oracle PL/SQL .......................................................................................... 1169
    41.12.1. Porting Examples ............................................................................................... 1170
    41.12.2. Other Things to Watch For ................................................................................... 1175
      41.12.2.1. Implicit Rollback after Exceptions ......................................................... 1175
      41.12.2.2. EXECUTE ............................................................................................. 1176
      41.12.2.3. Optimizing PL/pgSQL Functions ............................................................ 1176
  41.12.2.3. Optimizing PL/pgSQL Functions ........................................................................ 1176
CREATE USER MAPPING ................................................................. 1555
CREATE VIEW ................................................................. 1557
DEALLOCATE ............................................................... 1562
DECLARE ................................................................. 1563
DELETE ................................................................. 1567
DISCARD ................................................................. 1570
DO ................................................................. 1572
DROP ACCESS METHOD .......................................................... 1574
DROP AGGREGATE .......................................................... 1575
DROP CAST .............................................................. 1577
DROP COLLATION .......................................................... 1579
DROP CONVERSION ......................................................... 1581
DROP DATABASE ........................................................... 1582
DROP DOMAIN ............................................................ 1583
DROP EVENT TRIGGER ...................................................... 1584
DROP EXTENSION .......................................................... 1585
DROP FOREIGN DATA WRAPPER ........................................... 1587
DROP FOREIGN TABLE ...................................................... 1589
DROP FUNCTION .......................................................... 1591
DROP GROUP .............................................................. 1593
DROP INDEX ............................................................... 1594
DROP LANGUAGE .......................................................... 1596
DROP MATERIALIZED VIEW ................................................ 1598
DROP OPERATOR .......................................................... 1600
DROP OPERATOR CLASS .................................................. 1602
DROP OPERATOR FAMILY .................................................. 1604
DROP OWNED ............................................................. 1606
DROP POLICY ............................................................. 1608
DROP ROLE ................................................................. 1610
DROP RULE ................................................................. 1612
DROP SCHEMA ............................................................. 1614
DROP SEQUENCE .......................................................... 1616
DROP SERVER ............................................................. 1618
DROP TABLE ............................................................... 1620
DROP Tablespace ............................................................ 1622
DROP TEXT SEARCH CONFIGURATION .................................... 1624
DROP TEXT SEARCH DICTIONARY ......................................... 1626
DROP TEXT SEARCH PARSER ............................................... 1628
DROP TEXT SEARCH TEMPLATE ........................................... 1630
DROP TRANSFORM ........................................................ 1632
DROP TRIGGER ............................................................ 1634
DROP TYPE ................................................................. 1636
DROP USER ................................................................. 1638
DROP USER MAPPING ....................................................... 1639
DROP VIEW ................................................................. 1641
END ................................................................. 1642
EXECUTE ................................................................. 1643
EXPLAIN ................................................................. 1645
FETCH ................................................................. 1650
GRANT ................................................................. 1654
IMPORT FOREIGN SCHEMA ............................................... 1661
INSERT ................................................................. 1663
III. PostgreSQL Server Applications ................................................................. 1896
   initdb.................................................................................................................. 1897
   pg_archivecleanup .......................................................................................... 1901
   pg_controldata ................................................................................................. 1903
   pg_ctl ................................................................................................................. 1904
   pg_rewind ......................................................................................................... 1910
   pg_test_fsm ....................................................................................................... 1913
   pg_test_timing ................................................................................................. 1916
   pg_upgrade ........................................................................................................ 1922
   pg_xlogdump ...................................................................................................... 1929
   postgres ............................................................................................................. 1932
   postmaster ......................................................................................................... 1940

VII. Internals ........................................................................................................ 1941
  49. Overview of PostgreSQL Internals .............................................................. 1943
     49.1. The Path of a Query.................................................................................. 1943
     49.2. How Connections are Established.......................................................... 1943
     49.3. The Parser Stage...................................................................................... 1944
        49.3.1. Parser............................................................................................... 1944
        49.3.2. Transformation Process................................................................. 1945
     49.4. The PostgreSQL Rule System............................................................... 1945
     49.5. Planner/Optimizer................................................................................... 1945
        49.5.1. Generating Possible Plans............................................................... 1946
     49.6. Executor.................................................................................................. 1947
  50. System Catalogs............................................................................................ 1949
     50.1. Overview.................................................................................................. 1949
     50.2. pg_aggregate............................................................................................ 1950
     50.3. pg_am ....................................................................................................... 1953
     50.4. pg_amop ................................................................................................... 1954
     50.5. pg_amproc............................................................................................... 1955
     50.6. pg_attrdef ............................................................................................... 1955
     50.7. pg_attribute............................................................................................. 1956
     50.8. pg_authid................................................................................................ 1959
     50.9. pg_auth_members..................................................................................... 1960
     50.10. pg_cast .................................................................................................... 1961
     50.11. pg_class.................................................................................................. 1962
     50.12. pg_collation.......................................................................................... 1967
     50.13. pg_constraint........................................................................................ 1967
     50.14. pg_conversion....................................................................................... 1970
     50.15. pg_database ........................................................................................... 1971
     50.16. pg_db_role_setting................................................................................. 1973
     50.17. pg_default_acl....................................................................................... 1974
     50.18. pg_depend ................................................................................................ 1974
     50.19. pg_description...................................................................................... 1976
     50.20. pg_enum .................................................................................................. 1977
     50.21. pg_event_trigger................................................................................... 1977
     50.22. pg_extension.......................................................................................... 1978
     50.23. pg_foreign_data_wrapper................................................................. 1979
     50.24. pg_foreign_server................................................................................ 1980
     50.25. pg_foreign_table.................................................................................... 1981
61.4.1. GiST buffering build ................................................................. 2158
61.5. Examples ......................................................................................... 2159
62. SP-GiST Indexes .................................................................................. 2160
62.1. Introduction ...................................................................................... 2160
62.2. Built-in Operator Classes ................................................................. 2160
62.3. Extensibility ...................................................................................... 2160
62.4. Implementation ................................................................................ 2167
62.4.1. SP-GiST Limits ............................................................................. 2167
62.4.2. SP-GiST Without Node Labels ..................................................... 2168
62.4.3. “All-the-same” Inner Tuples .......................................................... 2168
62.5. Examples ......................................................................................... 2168
63. GIN Indexes ......................................................................................... 2169
63.1. Introduction ...................................................................................... 2169
63.2. Built-in Operator Classes ................................................................. 2169
63.3. Extensibility ...................................................................................... 2170
63.4. Implementation ................................................................................ 2173
63.4.1. GIN Fast Update Technique ......................................................... 2173
63.4.2. Partial Match Algorithm ............................................................... 2173
63.5. GIN Tips and Tricks ......................................................................... 2174
63.6. Limitations ....................................................................................... 2175
63.7. Examples ......................................................................................... 2175
64. BRIN Indexes ....................................................................................... 2176
64.1. Introduction ...................................................................................... 2176
64.1.1. Index Maintenance ...................................................................... 2176
64.2. Built-in Operator Classes ................................................................. 2176
64.3. Extensibility ...................................................................................... 2177
65. Database Physical Storage ................................................................... 2181
65.1. Database File Layout ....................................................................... 2181
65.2. TOAST ............................................................................................ 2183
65.2.1. Out-of-line, on-disk TOAST storage ............................................. 2184
65.2.2. Out-of-line, in-memory TOAST storage ....................................... 2185
65.3. Free Space Map ............................................................................... 2186
65.4. Visibility Map .................................................................................. 2186
65.5. The Initialization Fork ..................................................................... 2187
65.6. Database Page Layout ..................................................................... 2187
66. BKI Backend Interface ....................................................................... 2191
66.1. BKI File Format ............................................................................... 2191
66.2. BKI Commands ............................................................................... 2191
66.3. Structure of the Bootstrap BKI File .................................................. 2192
66.4. Example .......................................................................................... 2193
67. How the Planner Uses Statistics ....................................................... 2194
67.1. Row Estimation Examples ............................................................... 2194
67.2. Planner Statistics and Security ......................................................... 2199
VIII. Appendices ..................................................................................... 2201
A. PostgreSQL Error Codes ..................................................................... 2202
B. Date/Time Support ............................................................................... 2210
B.1. Date/Time Input Interpretation ......................................................... 2210
B.2. Handling of Invalid or Ambiguous Timestamps ................................ 2211
B.3. Date/Time Key Words ..................................................................... 2212
B.4. Date/Time Configuration Files ......................................................... 2213
B.5. History of Units ............................................................................... 2214
C. SQL Key Words .......................................................... 2216
D. SQL Conformance .......................................................... 2239
  D.1. Supported Features ............................................... 2240
  D.2. Unsupported Features ......................................... 2256
E. Release Notes .............................................................. 2271
  E.1. Release 9.6.1 .......................................................... 2271
    E.1.1. Migration to Version 9.6.1 .............................. 2271
    E.1.2. Changes ......................................................... 2271
  E.2. Release 9.6.16 ...................................................... 2273
    E.2.1. Migration to Version 9.6.16.......................... 2273
    E.2.2. Changes ......................................................... 2274
  E.3. Release 9.6.15 ...................................................... 2278
    E.3.1. Migration to Version 9.6.15......................... 2278
    E.3.2. Changes ......................................................... 2278
    E.4.1. Migration to Version 9.6.14......................... 2280
    E.4.2. Changes ......................................................... 2280
  E.5. Release 9.6.13 ...................................................... 2281
    E.5.1. Migration to Version 9.6.13.......................... 2282
    E.5.2. Changes ......................................................... 2282
  E.6. Release 9.6.12 ...................................................... 2284
    E.6.1. Migration to Version 9.6.12......................... 2284
    E.6.2. Changes ......................................................... 2285
  E.7. Release 9.6.11 ...................................................... 2288
    E.7.1. Migration to Version 9.6.11.......................... 2288
    E.7.2. Changes ......................................................... 2288
  E.8. Release 9.6.10 ...................................................... 2292
    E.8.1. Migration to Version 9.6.10......................... 2292
    E.8.2. Changes ......................................................... 2292
  E.9. Release 9.6.9 ...................................................... 2295
    E.9.1. Migration to Version 9.6.9.......................... 2295
    E.9.2. Changes ......................................................... 2295
  E.10. Release 9.6.8 ...................................................... 2299
    E.10.1. Migration to Version 9.6.8....................... 2299
    E.10.2. Changes ......................................................... 2299
  E.11. Release 9.6.7 ...................................................... 2300
    E.11.1. Migration to Version 9.6.7....................... 2301
    E.11.2. Changes ......................................................... 2301
  E.12. Release 9.6.6 ...................................................... 2304
    E.12.1. Migration to Version 9.6.6....................... 2304
    E.12.2. Changes ......................................................... 2304
  E.13. Release 9.6.5 ...................................................... 2307
    E.13.1. Migration to Version 9.6.5....................... 2307
    E.13.2. Changes ......................................................... 2307
  E.14. Release 9.6.4 ...................................................... 2308
    E.14.1. Migration to Version 9.6.4....................... 2308
    E.14.2. Changes ......................................................... 2309
  E.15. Release 9.6.3 ...................................................... 2314
    E.15.1. Migration to Version 9.6.3....................... 2314
    E.15.2. Changes ......................................................... 2314
  E.16. Release 9.6.2 ...................................................... 2318
    E.16.1. Migration to Version 9.6.2....................... 2318
Preface

This book is the official documentation of PostgreSQL. It has been written by the PostgreSQL developers and other volunteers in parallel to the development of the PostgreSQL software. It describes all the functionality that the current version of PostgreSQL officially supports.

To make the large amount of information about PostgreSQL manageable, this book has been organized in several parts. Each part is targeted at a different class of users, or at users in different stages of their PostgreSQL experience:

- Part I is an informal introduction for new users.
- Part II documents the SQL query language environment, including data types and functions, as well as user-level performance tuning. Every PostgreSQL user should read this.
- Part III describes the installation and administration of the server. Everyone who runs a PostgreSQL server, be it for private use or for others, should read this part.
- Part IV describes the programming interfaces for PostgreSQL client programs.
- Part V contains information for advanced users about the extensibility capabilities of the server. Topics include user-defined data types and functions.
- Part VI contains reference information about SQL commands, client and server programs. This part supports the other parts with structured information sorted by command or program.
- Part VII contains assorted information that might be of use to PostgreSQL developers.

1. What is PostgreSQL?

PostgreSQL is an object-relational database management system (ORDBMS) based on POSTGRES, Version 4.2\(^1\), developed at the University of California at Berkeley Computer Science Department. POSTGRES pioneered many concepts that only became available in some commercial database systems much later.

PostgreSQL is an open-source descendant of this original Berkeley code. It supports a large part of the SQL standard and offers many modern features:

- complex queries
- foreign keys
- triggers
- updatable views
- transactional integrity
- multiversion concurrency control

Also, PostgreSQL can be extended by the user in many ways, for example by adding new

- data types
- functions
- operators
- aggregate functions
- index methods

\(^1\) http://db.cs.berkeley.edu/postgres.html
And because of the liberal license, PostgreSQL can be used, modified, and distributed by anyone free of charge for any purpose, be it private, commercial, or academic.

2. A Brief History of PostgreSQL

The object-relational database management system now known as PostgreSQL is derived from the POSTGRES package written at the University of California at Berkeley. With over two decades of development behind it, PostgreSQL is now the most advanced open-source database available anywhere.

2.1. The Berkeley POSTGRES Project

The POSTGRES project, led by Professor Michael Stonebraker, was sponsored by the Defense Advanced Research Projects Agency (DARPA), the Army Research Office (ARO), the National Science Foundation (NSF), and ESL, Inc. The implementation of POSTGRES began in 1986. The initial concepts for the system were presented in The design of POSTGRES, and the definition of the initial data model appeared in The POSTGRES data model. The design of the rule system at that time was described in The design of the POSTGRES rules system. The rationale and architecture of the storage manager were detailed in The design of the POSTGRES storage system.

POSTGRES has undergone several major releases since then. The first “demoware” system became operational in 1987 and was shown at the 1988 ACM-SIGMOD Conference. Version 1, described in The implementation of POSTGRES, was released to a few external users in June 1989. In response to a critique of the first rule system (A commentary on the POSTGRES rules system), the rule system was redesigned (On Rules, Procedures, Caching and Views in Database Systems), and Version 2 was released in June 1990 with the new rule system. Version 3 appeared in 1991 and added support for multiple storage managers, an improved query executor, and a rewritten rule system. For the most part, subsequent releases until Postgres95 (see below) focused on portability and reliability.

POSTGRES has been used to implement many different research and production applications. These include: a financial data analysis system, a jet engine performance monitoring package, an asteroid tracking database, a medical information database, and several geographic information systems. POSTGRES has also been used as an educational tool at several universities. Finally, Illustra Information Technologies (later merged into Informix, which is now owned by IBM) picked up the code and commercialized it. In late 1992, POSTGRES became the primary data manager for the Sequoia 2000 scientific computing project.

The size of the external user community nearly doubled during 1993. It became increasingly obvious that maintenance of the prototype code and support was taking up large amounts of time that should have been devoted to database research. In an effort to reduce this support burden, the Berkeley POSTGRES project officially ended with Version 4.2.

4. http://meteora.ucsd.edu/s2k/s2k_home.html
2.2. Postgres95

In 1994, Andrew Yu and Jolly Chen added an SQL language interpreter to POSTGRES. Under a new name, Postgres95 was subsequently released to the web to find its own way in the world as an open-source descendant of the original POSTGRES Berkeley code.

Postgres95 code was completely ANSI C and trimmed in size by 25%. Many internal changes improved performance and maintainability. Postgres95 release 1.0.x ran about 30-50% faster on the Wisconsin Benchmark compared to POSTGRES, Version 4.2. Apart from bug fixes, the following were the major enhancements:

- The query language PostQUEL was replaced with SQL (implemented in the server). (Interface library libpq was named after PostQUEL.) Subqueries were not supported until PostgreSQL (see below), but they could be imitated in Postgres95 with user-defined SQL functions. Aggregate functions were re-implemented. Support for the GROUP BY query clause was also added.
- A new program (psql) was provided for interactive SQL queries, which used GNU Readline. This largely superseded the old monitor program.
- A new front-end library, libpgtcl, supported Tcl-based clients. A sample shell, pgtclsh, provided new Tcl commands to interface Tcl programs with the Postgres95 server.
- The large-object interface was overhauled. The inversion large objects were the only mechanism for storing large objects. (The inversion file system was removed.)
- The instance-level rule system was removed. Rules were still available as rewrite rules.
- A short tutorial introducing regular SQL features as well as those of Postgres95 was distributed with the source code.
- GNU make (instead of BSD make) was used for the build. Also, Postgres95 could be compiled with an unpatched GCC (data alignment of doubles was fixed).

2.3. PostgreSQL

By 1996, it became clear that the name “Postgres95” would not stand the test of time. We chose a new name, PostgreSQL, to reflect the relationship between the original POSTGRES and the more recent versions with SQL capability. At the same time, we set the version numbering to start at 6.0, putting the numbers back into the sequence originally begun by the Berkeley POSTGRES project.

Many people continue to refer to PostgreSQL as “Postgres” (now rarely in all capital letters) because of tradition or because it is easier to pronounce. This usage is widely accepted as a nickname or alias.

The emphasis during development of Postgres95 was on identifying and understanding existing problems in the server code. With PostgreSQL, the emphasis has shifted to augmenting features and capabilities, although work continues in all areas.

Details about what has happened in PostgreSQL since then can be found in Appendix E.

3. Conventions

The following conventions are used in the synopsis of a command: brackets ([ and ]) indicate optional parts. (In the synopsis of a Tcl command, question marks (?) are used instead, as is usual in Tcl.)
Braces ( { and } ) and vertical lines ( | ) indicate that you must choose one alternative. Dots ( . . . ) mean that the preceding element can be repeated.

Where it enhances the clarity, SQL commands are preceded by the prompt =>, and shell commands are preceded by the prompt $. Normally, prompts are not shown, though.

An administrator is generally a person who is in charge of installing and running the server. A user could be anyone who is using, or wants to use, any part of the PostgreSQL system. These terms should not be interpreted too narrowly; this book does not have fixed presumptions about system administration procedures.

4. Further Information

Besides the documentation, that is, this book, there are other resources about PostgreSQL:

Wiki
The PostgreSQL wiki contains the project’s FAQ (Frequently Asked Questions) list, TODO list, and detailed information about many more topics.

Web Site
The PostgreSQL web site carries details on the latest release and other information to make your work or play with PostgreSQL more productive.

Mailing Lists
The mailing lists are a good place to have your questions answered, to share experiences with other users, and to contact the developers. Consult the PostgreSQL web site for details.

Yourself!
PostgreSQL is an open-source project. As such, it depends on the user community for ongoing support. As you begin to use PostgreSQL, you will rely on others for help, either through the documentation or through the mailing lists. Consider contributing your knowledge back. Read the mailing lists and answer questions. If you learn something which is not in the documentation, write it up and contribute it. If you add features to the code, contribute them.

5. Bug Reporting Guidelines

When you find a bug in PostgreSQL we want to hear about it. Your bug reports play an important part in making PostgreSQL more reliable because even the utmost care cannot guarantee that every part of PostgreSQL will work on every platform under every circumstance.

The following suggestions are intended to assist you in forming bug reports that can be handled in an effective fashion. No one is required to follow them but doing so tends to be to everyone’s advantage.

We cannot promise to fix every bug right away. If the bug is obvious, critical, or affects a lot of users, chances are good that someone will look into it. It could also happen that we tell you to update to a newer version to see if the bug happens there. Or we might decide that the bug cannot be fixed before

5. https://wiki.postgresql.org
8. https://www.postgresql.org
some major rewrite we might be planning is done. Or perhaps it is simply too hard and there are more important things on the agenda. If you need help immediately, consider obtaining a commercial support contract.

5.1. Identifying Bugs

Before you report a bug, please read and re-read the documentation to verify that you can really do whatever it is you are trying. If it is not clear from the documentation whether you can do something or not, please report that too; it is a bug in the documentation. If it turns out that a program does something different from what the documentation says, that is a bug. That might include, but is not limited to, the following circumstances:

- A program terminates with a fatal signal or an operating system error message that would point to a problem in the program. (A counterexample might be a “disk full” message, since you have to fix that yourself.)
- A program produces the wrong output for any given input.
- A program refuses to accept valid input (as defined in the documentation).
- A program accepts invalid input without a notice or error message. But keep in mind that your idea of invalid input might be our idea of an extension or compatibility with traditional practice.
- PostgreSQL fails to compile, build, or install according to the instructions on supported platforms.

Here “program” refers to any executable, not only the backend process.

Being slow or resource-hogging is not necessarily a bug. Read the documentation or ask on one of the mailing lists for help in tuning your applications. Failing to comply to the SQL standard is not necessarily a bug either, unless compliance for the specific feature is explicitly claimed.

Before you continue, check on the TODO list and in the FAQ to see if your bug is already known. If you cannot decode the information on the TODO list, report your problem. The least we can do is make the TODO list clearer.

5.2. What to Report

The most important thing to remember about bug reporting is to state all the facts and only facts. Do not speculate what you think went wrong, what “it seemed to do”, or which part of the program has a fault. If you are not familiar with the implementation you would probably guess wrong and not help us a bit. And even if you are, educated explanations are a great supplement to but no substitute for facts. If we are going to fix the bug we still have to see it happen for ourselves first. Reporting the bare facts is relatively straightforward (you can probably copy and paste them from the screen) but all too often important details are left out because someone thought it does not matter or the report would be understood anyway.

The following items should be contained in every bug report:

- The exact sequence of steps from program start-up necessary to reproduce the problem. This should be self-contained; it is not enough to send in a bare SELECT statement without the preceding CREATE TABLE and INSERT statements, if the output should depend on the data in the tables. We do not have the time to reverse-engineer your database schema, and if we are supposed to make up our own data we would probably miss the problem.
The best format for a test case for SQL-related problems is a file that can be run through the psql frontend that shows the problem. (Be sure to not have anything in your ~/.psqlrc start-up file.) An easy way to create this file is to use pg_dump to dump out the table declarations and data needed to set the scene, then add the problem query. You are encouraged to minimize the size of your example, but this is not absolutely necessary. If the bug is reproducible, we will find it either way.

If your application uses some other client interface, such as PHP, then please try to isolate the offending queries. We will probably not set up a web server to reproduce your problem. In any case remember to provide the exact input files; do not guess that the problem happens for “large files” or “midsize databases”, etc. since this information is too inexact to be of use.

• The output you got. Please do not say that it “didn’t work” or “crashed”. If there is an error message, show it, even if you do not understand it. If the program terminates with an operating system error, say which. If nothing at all happens, say so. Even if the result of your test case is a program crash or otherwise obvious it might not happen on our platform. The easiest thing is to copy the output from the terminal, if possible.

**Note:** If you are reporting an error message, please obtain the most verbose form of the message. In psql, say \set VERBOSITY verbose beforehand. If you are extracting the message from the server log, set the run-time parameter log_error_verbosity to verbose so that all details are logged.

**Note:** In case of fatal errors, the error message reported by the client might not contain all the information available. Please also look at the log output of the database server. If you do not keep your server’s log output, this would be a good time to start doing so.

• The output you expected is very important to state. If you just write “This command gives me that output.” or “This is not what I expected.”, we might run it ourselves, scan the output, and think it looks OK and is exactly what we expected. We should not have to spend the time to decode the exact semantics behind your commands. Especially refrain from merely saying that “This is not what SQL says/Oracle does.” Digging out the correct behavior from SQL is not a fun undertaking, nor do we all know how all the other relational databases out there behave. (If your problem is a program crash, you can obviously omit this item.)

• Any command line options and other start-up options, including any relevant environment variables or configuration files that you changed from the default. Again, please provide exact information. If you are using a prepackaged distribution that starts the database server at boot time, you should try to find out how that is done.

• Anything you did at all differently from the installation instructions.

• The PostgreSQL version. You can run the command SELECT version(); to find out the version of the server you are connected to. Most executable programs also support a --version option; at least postgres --version and psql --version should work. If the function or the options do not exist then your version is more than old enough to warrant an upgrade. If you run a prepackaged version, such as RPMs, say so, including any subversion the package might have. If you are talking about a Git snapshot, mention that, including the commit hash.

If your version is older than 9.6.17 we will almost certainly tell you to upgrade. There are many bug fixes and improvements in each new release, so it is quite possible that a bug you have encountered
in an older release of PostgreSQL has already been fixed. We can only provide limited support
for sites using older releases of PostgreSQL; if you require more than we can provide, consider
acquiring a commercial support contract.

- Platform information. This includes the kernel name and version, C library, processor, memory
  information, and so on. In most cases it is sufficient to report the vendor and version, but do not
  assume everyone knows what exactly “Debian” contains or that everyone runs on x86_64. If you
  have installation problems then information about the toolchain on your machine (compiler, make,
  and so on) is also necessary.

Do not be afraid if your bug report becomes rather lengthy. That is a fact of life. It is better to report
everything the first time than us having to squeeze the facts out of you. On the other hand, if your
input files are huge, it is fair to ask first whether somebody is interested in looking into it. Here is an
article\footnote{http://www.chiark.greenend.org.uk/~sgtatham/bugs.html} that outlines some more tips on reporting bugs.

Do not spend all your time to figure out which changes in the input make the problem go away. This
will probably not help solving it. If it turns out that the bug cannot be fixed right away, you will still
have time to find and share your work-around. Also, once again, do not waste your time guessing why
the bug exists. We will find that out soon enough.

When writing a bug report, please avoid confusing terminology. The software package in total is
called “PostgreSQL”, sometimes “Postgres” for short. If you are specifically talking about the back-
end process, mention that, do not just say “PostgreSQL crashes”. A crash of a single backend process
is quite different from crash of the parent “postgres” process; please don’t say “the server crashed”
when you mean a single backend process went down, nor vice versa. Also, client programs such as the
interactive frontend “psql” are completely separate from the backend. Please try to be specific about
whether the problem is on the client or server side.

5.3. Where to Report Bugs

In general, send bug reports to the bug report mailing list at
\texttt{<pgsql-bugs@lists.postgresql.org>}. You are requested to use a
descriptive subject for your email message, perhaps parts of the error message.

Another method is to fill in the bug report web-form available at the project’s web site\footnote{https://www.postgresql.org/}. Entering a
bug report this way causes it to be mailed to the \texttt{<pgsql-bugs@lists.postgresql.org>} mailing
list.

If your bug report has security implications and you’d prefer that it not become immediately vis-
ible in public archives, don’t send it to \texttt{pgsql-bugs}. Security issues can be reported privately to
\texttt{<security@postgresql.org>}.

Do not send bug reports to any of the user mailing lists, such as
\texttt{<pgsql-sql@lists.postgresql.org>} or \texttt{<pgsql-general@lists.postgresql.org>}. These mailing lists are for answering user questions, and their subscribers normally do not wish to
receive bug reports. More importantly, they are unlikely to fix them.

Also, please do \textbf{not} send reports to the developers’ mailing list
\texttt{<pgsql-hackers@lists.postgresql.org>}. This list is for discussing the development of
PostgreSQL, and it would be nice if we could keep the bug reports separate. We might choose to take
up a discussion about your bug report on \texttt{pgsql-hackers}, if the problem needs more review.

\footnote{http://www.chiark.greenend.org.uk/~sgtatham/bugs.html}
\footnote{https://www.postgresql.org/}
If you have a problem with the documentation, the best place to report it is the documentation mailing list <pgsql-docs@lists.postgresql.org>. Please be specific about what part of the documentation you are unhappy with.

If your bug is a portability problem on a non-supported platform, send mail to <pgsql-hackers@lists.postgresql.org>, so we (and you) can work on porting PostgreSQL to your platform.

**Note:** Due to the unfortunate amount of spam going around, all of the above lists will be moderated unless you are subscribed. That means there will be some delay before the email is delivered. If you wish to subscribe to the lists, please visit https://lists.postgresql.org/ for instructions.
I. Tutorial

Welcome to the PostgreSQL Tutorial. The following few chapters are intended to give a simple introduction to PostgreSQL, relational database concepts, and the SQL language to those who are new to any one of these aspects. We only assume some general knowledge about how to use computers. No particular Unix or programming experience is required. This part is mainly intended to give you some hands-on experience with important aspects of the PostgreSQL system. It makes no attempt to be a complete or thorough treatment of the topics it covers.

After you have worked through this tutorial you might want to move on to reading Part II to gain a more formal knowledge of the SQL language, or Part IV for information about developing applications for PostgreSQL. Those who set up and manage their own server should also read Part III.
Chapter 1. Getting Started

1.1. Installation

Before you can use PostgreSQL you need to install it, of course. It is possible that PostgreSQL is already installed at your site, either because it was included in your operating system distribution or because the system administrator already installed it. If that is the case, you should obtain information from the operating system documentation or your system administrator about how to access PostgreSQL.

If you are not sure whether PostgreSQL is already available or whether you can use it for your experimentation then you can install it yourself. Doing so is not hard and it can be a good exercise. PostgreSQL can be installed by any unprivileged user; no superuser (root) access is required.

If you are installing PostgreSQL yourself, then refer to Chapter 16 for instructions on installation, and return to this guide when the installation is complete. Be sure to follow closely the section about setting up the appropriate environment variables.

If your site administrator has not set things up in the default way, you might have some more work to do. For example, if the database server machine is a remote machine, you will need to set the `PGHOST` environment variable to the name of the database server machine. The environment variable `PGPORT` might also have to be set. The bottom line is this: if you try to start an application program and it complains that it cannot connect to the database, you should consult your site administrator or, if that is you, the documentation to make sure that your environment is properly set up. If you did not understand the preceding paragraph then read the next section.

1.2. Architectural Fundamentals

Before we proceed, you should understand the basic PostgreSQL system architecture. Understanding how the parts of PostgreSQL interact will make this chapter somewhat clearer.

In database jargon, PostgreSQL uses a client/server model. A PostgreSQL session consists of the following cooperating processes (programs):

- A server process, which manages the database files, accepts connections to the database from client applications, and performs database actions on behalf of the clients. The database server program is called `postgres`.
- The user’s client (frontend) application that wants to perform database operations. Client applications can be very diverse in nature: a client could be a text-oriented tool, a graphical application, a web server that accesses the database to display web pages, or a specialized database maintenance tool. Some client applications are supplied with the PostgreSQL distribution; most are developed by users.

As is typical of client/server applications, the client and the server can be on different hosts. In that case they communicate over a TCP/IP network connection. You should keep this in mind, because the files that can be accessed on a client machine might not be accessible (or might only be accessible using a different file name) on the database server machine.

The PostgreSQL server can handle multiple concurrent connections from clients. To achieve this it starts (“forks”) a new process for each connection. From that point on, the client and the new
server process communicate without intervention by the original \texttt{postgres} process. Thus, the master server process is always running, waiting for client connections, whereas client and associated server processes come and go. (All of this is of course invisible to the user. We only mention it here for completeness.)

1.3. Creating a Database

The first test to see whether you can access the database server is to try to create a database. A running PostgreSQL server can manage many databases. Typically, a separate database is used for each project or for each user.

Possibly, your site administrator has already created a database for your use. In that case you can omit this step and skip ahead to the next section.

To create a new database, in this example named \texttt{mydb}, you use the following command:

\begin{verbatim}
$ createdb mydb
\end{verbatim}

If this produces no response then this step was successful and you can skip over the remainder of this section.

If you see a message similar to:

\begin{verbatim}
createdb: command not found
\end{verbatim}

then PostgreSQL was not installed properly. Either it was not installed at all or your shell’s search path was not set to include it. Try calling the command with an absolute path instead:

\begin{verbatim}
$ /usr/local/pgsql/bin/createdb mydb
\end{verbatim}

The path at your site might be different. Contact your site administrator or check the installation instructions to correct the situation.

Another response could be this:

\begin{verbatim}
createdb: could not connect to database postgres: could not connect to server: No such file or directory
Is the server running locally and accepting connections on Unix domain socket "/tmp/.s.PGSQL.5432"?
\end{verbatim}

This means that the server was not started, or it was not started where \texttt{createdb} expected it. Again, check the installation instructions or consult the administrator.

Another response could be this:

\begin{verbatim}
createdb: could not connect to database postgres: FATAL: role "joe" does not exist
\end{verbatim}

where your own login name is mentioned. This will happen if the administrator has not created a PostgreSQL user account for you. (PostgreSQL user accounts are distinct from operating system user accounts.) If you are the administrator, see Chapter 21 for help creating accounts. You will need to become the operating system user under which PostgreSQL was installed (usually \texttt{postgres}) to create the first user account. It could also be that you were assigned a PostgreSQL user name that is different from your operating system user name; in that case you need to use the \texttt{-U} switch or set the \texttt{PGUSER} environment variable to specify your PostgreSQL user name.

If you have a user account but it does not have the privileges required to create a database, you will see the following:
Chapter 1. Getting Started

createdb: database creation failed: ERROR: permission denied to create database

Not every user has authorization to create new databases. If PostgreSQL refuses to create databases for you then the site administrator needs to grant you permission to create databases. Consult your site administrator if this occurs. If you installed PostgreSQL yourself then you should log in for the purposes of this tutorial under the user account that you started the server as. ¹

You can also create databases with other names. PostgreSQL allows you to create any number of databases at a given site. Database names must have an alphabetic first character and are limited to 63 bytes in length. A convenient choice is to create a database with the same name as your current user name. Many tools assume that database name as the default, so it can save you some typing. To create that database, simply type:

$ createdb

If you do not want to use your database anymore you can remove it. For example, if you are the owner (creator) of the database mydb, you can destroy it using the following command:

$ dropdb mydb

(For this command, the database name does not default to the user account name. You always need to specify it.) This action physically removes all files associated with the database and cannot be undone, so this should only be done with a great deal of forethought.

More about createdb and dropdb can be found in createdb and dropdb respectively.

1.4. Accessing a Database

Once you have created a database, you can access it by:

- Running the PostgreSQL interactive terminal program, called psql, which allows you to interactively enter, edit, and execute SQL commands.
- Using an existing graphical frontend tool like pgAdmin or an office suite with ODBC or JDBC support to create and manipulate a database. These possibilities are not covered in this tutorial.
- Writing a custom application, using one of the several available language bindings. These possibilities are discussed further in Part IV.

You probably want to start up psql to try the examples in this tutorial. It can be activated for the mydb database by typing the command:

$ psql mydb

If you do not supply the database name then it will default to your user account name. You already discovered this scheme in the previous section using createdb.

In psql, you will be greeted with the following message:

¹ As an explanation for why this works: PostgreSQL user names are separate from operating system user accounts. When you connect to a database, you can choose what PostgreSQL user name to connect as; if you don’t, it will default to the same name as your current operating system account. As it happens, there will always be a PostgreSQL user account that has the same name as the operating system user that started the server, and it also happens that that user always has permission to create databases. Instead of logging in as that user you can also specify the –U option everywhere to select a PostgreSQL user name to connect as.
psql (9.6.17)
Type "help" for help.

mydb=>

The last line could also be:

mydb=#

That would mean you are a database superuser, which is most likely the case if you installed the PostgreSQL instance yourself. Being a superuser means that you are not subject to access controls. For the purposes of this tutorial that is not important.

If you encounter problems starting psql then go back to the previous section. The diagnostics of createdb and psql are similar, and if the former worked the latter should work as well.

The last line printed out by psql is the prompt, and it indicates that psql is listening to you and that you can type SQL queries into a work space maintained by psql. Try out these commands:

mydb=> SELECT version();

<table>
<thead>
<tr>
<th>version</th>
</tr>
</thead>
<tbody>
<tr>
<td>PostgreSQL 9.6.17 on x86_64-pc-linux-gnu, compiled by gcc (Debian 4.9.2-10) 4.9.2, 64-bit</td>
</tr>
</tbody>
</table>

mydb=> SELECT current_date;

date
| 2016-01-07 |

(1 row)

mydb=> SELECT 2 + 2;

<table>
<thead>
<tr>
<th>?column?</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
</tr>
</tbody>
</table>

(1 row)

The psql program has a number of internal commands that are not SQL commands. They begin with the backslash character, "\". For example, you can get help on the syntax of various PostgreSQL SQL commands by typing:

mydb=> \h

To get out of psql, type:

mydb=> \q

and psql will quit and return you to your command shell. (For more internal commands, type \? at the psql prompt.) The full capabilities of psql are documented in psql. In this tutorial we will not use these features explicitly, but you can use them yourself when it is helpful.
Chapter 2. The SQL Language

2.1. Introduction

This chapter provides an overview of how to use SQL to perform simple operations. This tutorial is only intended to give you an introduction and is in no way a complete tutorial on SQL. Numerous books have been written on SQL, including Understanding the New SQL and A Guide to the SQL Standard. You should be aware that some PostgreSQL language features are extensions to the standard.

In the examples that follow, we assume that you have created a database named mydb, as described in the previous chapter, and have been able to start psql.

Examples in this manual can also be found in the PostgreSQL source distribution in the directory src/tutorial/. (Binary distributions of PostgreSQL might not compile these files.) To use those files, first change to that directory and run make:

```
$ cd ....../src/tutorial
$ make
```

This creates the scripts and compiles the C files containing user-defined functions and types. Then, to start the tutorial, do the following:

```
$ cd ....../tutorial
$ psql -s mydb
...
```

```
mydb=> \i basics.sql
```

The \i command reads in commands from the specified file. psql's -s option puts you in single step mode which pauses before sending each statement to the server. The commands used in this section are in the file basics.sql.

2.2. Concepts

PostgreSQL is a relational database management system (RDBMS). That means it is a system for managing data stored in relations. Relation is essentially a mathematical term for table. The notion of storing data in tables is so commonplace today that it might seem inherently obvious, but there are a number of other ways of organizing databases. Files and directories on Unix-like operating systems form an example of a hierarchical database. A more modern development is the object-oriented database.

Each table is a named collection of rows. Each row of a given table has the same set of named columns, and each column is of a specific data type. Whereas columns have a fixed order in each row, it is important to remember that SQL does not guarantee the order of the rows within the table in any way (although they can be explicitly sorted for display).

Tables are grouped into databases, and a collection of databases managed by a single PostgreSQL server instance constitutes a database cluster.
2.3. Creating a New Table

You can create a new table by specifying the table name, along with all column names and their types:

```
CREATE TABLE weather (  
city varchar(80),       
temp_lo int, -- low temperature  
temp_hi int, -- high temperature  
prcp real, -- precipitation  
date date  
);  
```

You can enter this into `psql` with the line breaks. `psql` will recognize that the command is not terminated until the semicolon.

White space (i.e., spaces, tabs, and newlines) can be used freely in SQL commands. That means you can type the command aligned differently than above, or even all on one line. Two dashes (“`--`”) introduce comments. Whatever follows them is ignored up to the end of the line. SQL is case insensitive about key words and identifiers, except when identifiers are double-quoted to preserve the case (not done above).

`varchar(80)` specifies a data type that can store arbitrary character strings up to 80 characters in length. `int` is the normal integer type. `real` is a type for storing single precision floating-point numbers. `date` should be self-explanatory. (Yes, the column of type `date` is also named `date`. This might be convenient or confusing — you choose.)

PostgreSQL supports the standard SQL types `int`, `smallint`, `real`, `double precision`, `char(N)`, `varchar(N)`, `date`, `time`, `timestamp`, and `interval`, as well as other types of general utility and a rich set of geometric types. PostgreSQL can be customized with an arbitrary number of user-defined data types. Consequently, type names are not key words in the syntax, except where required to support special cases in the SQL standard.

The second example will store cities and their associated geographical location:

```
CREATE TABLE cities (  
name varchar(80),  
location point  
);  
```

The `point` type is an example of a PostgreSQL-specific data type.

Finally, it should be mentioned that if you don’t need a table any longer or want to recreate it differently you can remove it using the following command:

```
DROP TABLE tablename;  
```

2.4. Populating a Table With Rows

The `INSERT` statement is used to populate a table with rows:

```
INSERT INTO weather VALUES ('San Francisco', 46, 50, 0.25, '1994-11-27');  
```
Note that all data types use rather obvious input formats. Constants that are not simple numeric values usually must be surrounded by single quotes (’), as in the example. The **date** type is actually quite flexible in what it accepts, but for this tutorial we will stick to the unambiguous format shown here.

The **point** type requires a coordinate pair as input, as shown here:

```sql
INSERT INTO cities VALUES ('San Francisco', '(-194.0, 53.0)');
```

The syntax used so far requires you to remember the order of the columns. An alternative syntax allows you to list the columns explicitly:

```sql
INSERT INTO weather (city, temp_lo, temp_hi, prcp, date)
VALUES ('San Francisco', 43, 57, 0.0, '1994-11-29');
```

You can list the columns in a different order if you wish or even omit some columns, e.g., if the precipitation is unknown:

```sql
INSERT INTO weather (date, city, temp_hi, temp_lo)
VALUES ('1994-11-29', 'Hayward', 54, 37);
```

Many developers consider explicitly listing the columns better style than relying on the order implicitly.

Please enter all the commands shown above so you have some data to work with in the following sections.

You could also have used **COPY** to load large amounts of data from flat-text files. This is usually faster because the **COPY** command is optimized for this application while allowing less flexibility than **INSERT**. An example would be:

```sql
COPY weather FROM '/home/user/weather.txt';
```

where the file name for the source file must be available on the machine running the backend process, not the client, since the backend process reads the file directly. You can read more about the **COPY** command in **COPY**.

### 2.5. Querying a Table

To retrieve data from a table, the table is **queried**. An SQL **SELECT** statement is used to do this. The statement is divided into a select list (the part that lists the columns to be returned), a table list (the part that lists the tables from which to retrieve the data), and an optional qualification (the part that specifies any restrictions). For example, to retrieve all the rows of table **weather**, type:

```sql
SELECT * FROM weather;
```

Here * is a shorthand for “all columns”.¹ So the same result would be had with:

```sql
SELECT city, temp_lo, temp_hi, prcp, date FROM weather;
```

The output should be:

```
city | temp_lo | temp_hi | prcp | date
```

¹ While **SELECT * is useful for off-the-cuff queries, it is widely considered bad style in production code, since adding a column to the table would change the results.
You can write expressions, not just simple column references, in the select list. For example, you can do:

```sql
SELECT city, (temp_hi+temp_lo)/2 AS temp_avg, date FROM weather;
```

This should give:

<table>
<thead>
<tr>
<th>city</th>
<th>temp_avg</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>48</td>
<td>1994-11-27</td>
</tr>
<tr>
<td>San Francisco</td>
<td>50</td>
<td>1994-11-29</td>
</tr>
<tr>
<td>Hayward</td>
<td>45</td>
<td>1994-11-29</td>
</tr>
</tbody>
</table>

(3 rows)

Notice how the AS clause is used to relabel the output column. (The AS clause is optional.)

A query can be “qualified” by adding a WHERE clause that specifies which rows are wanted. The WHERE clause contains a Boolean (truth value) expression, and only rows for which the Boolean expression is true are returned. The usual Boolean operators (AND, OR, and NOT) are allowed in the qualification. For example, the following retrieves the weather of San Francisco on rainy days:

```sql
SELECT * FROM weather
WHERE city = 'San Francisco' AND prcp > 0.0;
```

Result:

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
</tbody>
</table>

(1 row)

You can request that the results of a query be returned in sorted order:

```sql
SELECT * FROM weather
ORDER BY city;
```

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>Hayward</td>
<td>37</td>
<td>54</td>
<td></td>
<td>1994-11-29</td>
</tr>
<tr>
<td>San Francisco</td>
<td>43</td>
<td>57</td>
<td>0</td>
<td>1994-11-29</td>
</tr>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
</tbody>
</table>

In this example, the sort order isn’t fully specified, and so you might get the San Francisco rows in either order. But you’d always get the results shown above if you do:

```sql
SELECT * FROM weather
ORDER BY city, temp_lo;
```
You can request that duplicate rows be removed from the result of a query:

```sql
SELECT DISTINCT city
FROM weather;
```

<table>
<thead>
<tr>
<th>city</th>
</tr>
</thead>
<tbody>
<tr>
<td>Hayward</td>
</tr>
<tr>
<td>San Francisco</td>
</tr>
</tbody>
</table>

(2 rows)

Here again, the result row ordering might vary. You can ensure consistent results by using `DISTINCT` and `ORDER BY` together:

```sql
SELECT DISTINCT city
FROM weather
ORDER BY city;
```

### 2.6. Joins Between Tables

Thus far, our queries have only accessed one table at a time. Queries can access multiple tables at once, or access the same table in such a way that multiple rows of the table are being processed at the same time. A query that accesses multiple rows of the same or different tables at one time is called a *join* query. As an example, say you wish to list all the weather records together with the location of the associated city. To do that, we need to compare the `city` column of each row of the `weather` table with the `name` column of all rows in the `cities` table, and select the pairs of rows where these values match.

**Note:** This is only a conceptual model. The join is usually performed in a more efficient manner than actually comparing each possible pair of rows, but this is invisible to the user.

This would be accomplished by the following query:

```sql
SELECT *
FROM weather, cities
WHERE city = name;
```

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
<th>name</th>
<th>location</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
<td>San Francisco</td>
<td>(-194,53)</td>
</tr>
<tr>
<td>San Francisco</td>
<td>43</td>
<td>57</td>
<td>0</td>
<td>1994-11-29</td>
<td>San Francisco</td>
<td>(-194,53)</td>
</tr>
</tbody>
</table>

(2 rows)

Observe two things about the result set:

- There is no result row for the city of Hayward. This is because there is no matching entry in the `cities` table for Hayward, so the join ignores the unmatched rows in the `weather` table. We will see shortly how this can be fixed.

2. In some database systems, including older versions of PostgreSQL, the implementation of `DISTINCT` automatically orders the rows and so `ORDER BY` is unnecessary. But this is not required by the SQL standard, and current PostgreSQL does not guarantee that `DISTINCT` causes the rows to be ordered.
• There are two columns containing the city name. This is correct because the lists of columns from the weather and cities tables are concatenated. In practice this is undesirable, though, so you will probably want to list the output columns explicitly rather than using *:

```
SELECT city, temp_lo, temp_hi, prcp, date, location
FROM weather, cities
WHERE city = name;
```

**Exercise:** Attempt to determine the semantics of this query when the **WHERE** clause is omitted.

Since the columns all had different names, the parser automatically found which table they belong to. If there were duplicate column names in the two tables you’d need to qualify the column names to show which one you meant, as in:

```
SELECT weather.city, weather.temp_lo, weather.temp_hi, weather.prcp, weather.date, cities.location
FROM weather, cities
WHERE cities.name = weather.city;
```

It is widely considered good style to qualify all column names in a join query, so that the query won’t fail if a duplicate column name is later added to one of the tables.

Join queries of the kind seen thus far can also be written in this alternative form:

```
SELECT *
FROM weather INNER JOIN cities ON (weather.city = cities.name);
```

This syntax is not as commonly used as the one above, but we show it here to help you understand the following topics.

Now we will figure out how we can get the Hayward records back in. What we want the query to do is to scan the weather table and for each row to find the matching cities row(s). If no matching row is found we want some “empty values” to be substituted for the cities table’s columns. This kind of query is called an **outer join**. (The joins we have seen so far are inner joins.) The command looks like this:

```
SELECT *
FROM weather LEFT OUTER JOIN cities ON (weather.city = cities.name);
```

This query is called a **left outer join** because the table mentioned on the left of the join operator will have each of its rows in the output at least once, whereas the table on the right will only have those rows output that match some row of the left table. When outputting a left-table row for which there is no right-table match, empty (null) values are substituted for the right-table columns.

**Exercise:** There are also right outer joins and full outer joins. Try to find out what those do.

We can also join a table against itself. This is called a **self join**. As an example, suppose we wish to find all the weather records that are in the temperature range of other weather records. So we need to compare the **temp_lo** and **temp_hi** columns of each weather row to the **temp_lo** and **temp_hi** columns of all other weather rows. We can do this with the following query:
SELECT W1.city, W1.temp_lo AS low, W1.temp_hi AS high, 
    W2.city, W2.temp_lo AS low, W2.temp_hi AS high 
FROM weather W1, weather W2 
WHERE W1.temp_lo < W2.temp_lo 
AND W1.temp_hi > W2.temp_hi;

<table>
<thead>
<tr>
<th>city</th>
<th>low</th>
<th>high</th>
<th>city</th>
<th>low</th>
<th>high</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>43</td>
<td>57</td>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
</tr>
<tr>
<td>Hayward</td>
<td>37</td>
<td>54</td>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
</tr>
</tbody>
</table>

Here we have relabeled the weather table as W1 and W2 to be able to distinguish the left and right side of the join. You can also use these kinds of aliases in other queries to save some typing, e.g.:

```sql
SELECT *
FROM weather w, cities c
WHERE w.city = c.name;
```

You will encounter this style of abbreviating quite frequently.

### 2.7. Aggregate Functions

Like most other relational database products, PostgreSQL supports aggregate functions. An aggregate function computes a single result from multiple input rows. For example, there are aggregates to compute the **count**, **sum**, **avg** (average), **max** (maximum) and **min** (minimum) over a set of rows.

As an example, we can find the highest low-temperature reading anywhere with:

```sql
SELECT max(temp_lo) FROM weather;
```

```
max
----
46
```

(1 row)

If we wanted to know what city (or cities) that reading occurred in, we might try:

```sql
SELECT city FROM weather WHERE temp_lo = max(temp_lo);
```

```
WRONG
```

but this will not work since the aggregate `max` cannot be used in the `WHERE` clause. (This restriction exists because the `WHERE` clause determines which rows will be included in the aggregate calculation; so obviously it has to be evaluated before aggregate functions are computed.) However, as is often the case the query can be restated to accomplish the desired result, here by using a subquery:

```sql
SELECT city FROM weather
    WHERE temp_lo = (SELECT max(temp_lo) FROM weather);
```

```
<table>
<thead>
<tr>
<th>city</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
</tr>
</tbody>
</table>
```

(1 row)
This is OK because the subquery is an independent computation that computes its own aggregate separately from what is happening in the outer query.

Aggregates are also very useful in combination with GROUP BY clauses. For example, we can get the maximum low temperature observed in each city with:

```
SELECT city, max(temp_lo)
FROM weather
GROUP BY city;
```

<table>
<thead>
<tr>
<th>city</th>
<th>max</th>
</tr>
</thead>
<tbody>
<tr>
<td>Hayward</td>
<td>37</td>
</tr>
<tr>
<td>San Francisco</td>
<td>46</td>
</tr>
</tbody>
</table>

(2 rows)

which gives us one output row per city. Each aggregate result is computed over the table rows matching that city. We can filter these grouped rows using HAVING:

```
SELECT city, max(temp_lo)
FROM weather
GROUP BY city
HAVING max(temp_lo) < 40;
```

<table>
<thead>
<tr>
<th>city</th>
<th>max</th>
</tr>
</thead>
<tbody>
<tr>
<td>Hayward</td>
<td>37</td>
</tr>
</tbody>
</table>

(1 row)

which gives us the same results for only the cities that have all temp_lo values below 40. Finally, if we only care about cities whose names begin with “s”, we might do:

```
SELECT city, max(temp_lo)
FROM weather
WHERE city LIKE 'S%'
GROUP BY city
HAVING max(temp_lo) < 40;
```

1 The LIKE operator does pattern matching and is explained in Section 9.7.

It is important to understand the interaction between aggregates and SQL’s WHERE and HAVING clauses. The fundamental difference between WHERE and HAVING is this: WHERE selects input rows before groups and aggregates are computed (thus, it controls which rows go into the aggregate computation), whereas HAVING selects group rows after groups and aggregates are computed. Thus, the WHERE clause must not contain aggregate functions; it makes no sense to try to use an aggregate to determine which rows will be inputs to the aggregates. On the other hand, the HAVING clause always contains aggregate functions. (Strictly speaking, you are allowed to write a HAVING clause that doesn’t use aggregates, but it’s seldom useful. The same condition could be used more efficiently at the WHERE stage.)

In the previous example, we can apply the city name restriction in WHERE, since it needs no aggregate. This is more efficient than adding the restriction to HAVING, because we avoid doing the grouping and aggregate calculations for all rows that fail the WHERE check.
2.8. Updates

You can update existing rows using the `UPDATE` command. Suppose you discover the temperature readings are all off by 2 degrees after November 28. You can correct the data as follows:

```
UPDATE weather
    SET temp_hi = temp_hi - 2, temp_lo = temp_lo - 2
    WHERE date > '1994-11-28';
```

Look at the new state of the data:

```
SELECT * FROM weather;
```

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
<tr>
<td>San Francisco</td>
<td>41</td>
<td>55</td>
<td>0</td>
<td>1994-11-29</td>
</tr>
<tr>
<td>Hayward</td>
<td>35</td>
<td>52</td>
<td></td>
<td>1994-11-29</td>
</tr>
</tbody>
</table>
(3 rows)

2.9. Deletions

Rows can be removed from a table using the `DELETE` command. Suppose you are no longer interested in the weather of Hayward. Then you can do the following to delete those rows from the table:

```
DELETE FROM weather WHERE city = 'Hayward';
```

All weather records belonging to Hayward are removed.

```
SELECT * FROM weather;
```

<table>
<thead>
<tr>
<th>city</th>
<th>temp_lo</th>
<th>temp_hi</th>
<th>prcp</th>
<th>date</th>
</tr>
</thead>
<tbody>
<tr>
<td>San Francisco</td>
<td>46</td>
<td>50</td>
<td>0.25</td>
<td>1994-11-27</td>
</tr>
<tr>
<td>San Francisco</td>
<td>41</td>
<td>55</td>
<td>0</td>
<td>1994-11-29</td>
</tr>
</tbody>
</table>
(2 rows)

One should be wary of statements of the form

```
DELETE FROM tablename;
```

Without a qualification, `DELETE` will remove all rows from the given table, leaving it empty. The system will not request confirmation before doing this!
Chapter 3. Advanced Features

3.1. Introduction

In the previous chapter we have covered the basics of using SQL to store and access your data in PostgreSQL. We will now discuss some more advanced features of SQL that simplify management and prevent loss or corruption of your data. Finally, we will look at some PostgreSQL extensions.

This chapter will on occasion refer to examples found in Chapter 2 to change or improve them, so it will be useful to have read that chapter. Some examples from this chapter can also be found in advanced.sql in the tutorial directory. This file also contains some sample data to load, which is not repeated here. (Refer to Section 2.1 for how to use the file.)

3.2. Views

Refer back to the queries in Section 2.6. Suppose the combined listing of weather records and city location is of particular interest to your application, but you do not want to type the query each time you need it. You can create a view over the query, which gives a name to the query that you can refer to like an ordinary table:

```
CREATE VIEW myview AS
    SELECT city, temp_lo, temp_hi, prcp, date, location
    FROM weather, cities
    WHERE city = name;
```

```
SELECT * FROM myview;
```

Making liberal use of views is a key aspect of good SQL database design. Views allow you to encapsulate the details of the structure of your tables, which might change as your application evolves, behind consistent interfaces.

Views can be used in almost any place a real table can be used. Building views upon other views is not uncommon.

3.3. Foreign Keys

Recall the weather and cities tables from Chapter 2. Consider the following problem: You want to make sure that no one can insert rows in the weather table that do not have a matching entry in the cities table. This is called maintaining the referential integrity of your data. In simplistic database systems this would be implemented (if at all) by first looking at the cities table to check if a matching record exists, and then inserting or rejecting the new weather records. This approach has a number of problems and is very inconvenient, so PostgreSQL can do this for you.

The new declaration of the tables would look like this:

```
CREATE TABLE cities {
    city varchar(80) primary key,
    location point
};
```
CREATE TABLE weather (
    city varchar(80) references cities(city),
    temp_lo int,
    temp_hi int,
    prcp real,
    date date
);

Now try inserting an invalid record:

INSERT INTO weather VALUES ('Berkeley', 45, 53, 0.0, '1994-11-28');

ERROR: insert or update on table "weather" violates foreign key constraint "weather_city_fkey"
DETAIL: Key (city)=(Berkeley) is not present in table "cities".

The behavior of foreign keys can be finely tuned to your application. We will not go beyond this simple example in this tutorial, but just refer you to Chapter 5 for more information. Making correct use of foreign keys will definitely improve the quality of your database applications, so you are strongly encouraged to learn about them.

3.4. Transactions

Transactions are a fundamental concept of all database systems. The essential point of a transaction is that it bundles multiple steps into a single, all-or-nothing operation. The intermediate states between the steps are not visible to other concurrent transactions, and if some failure occurs that prevents the transaction from completing, then none of the steps affect the database at all.

For example, consider a bank database that contains balances for various customer accounts, as well as total deposit balances for branches. Suppose that we want to record a payment of $100.00 from Alice’s account to Bob’s account. Simplifying outrageously, the SQL commands for this might look like:

```sql
UPDATE accounts SET balance = balance - 100.00
    WHERE name = 'Alice';
UPDATE branches SET balance = balance - 100.00
    WHERE name = (SELECT branch_name FROM accounts WHERE name = 'Alice');
UPDATE accounts SET balance = balance + 100.00
    WHERE name = 'Bob';
UPDATE branches SET balance = balance + 100.00
    WHERE name = (SELECT branch_name FROM accounts WHERE name = 'Bob');
```

The details of these commands are not important here; the important point is that there are several separate updates involved to accomplish this rather simple operation. Our bank’s officers will want to be assured that either all these updates happen, or none of them happen. It would certainly not do for a system failure to result in Bob receiving $100.00 that was not debited from Alice. Nor would Alice long remain a happy customer if she was debited without Bob being credited. We need a guarantee that if something goes wrong partway through the operation, none of the steps executed so far will take effect. Grouping the updates into a transaction gives us this guarantee. A transaction is said to be atomic: from the point of view of other transactions, it either happens completely or not at all.
We also want a guarantee that once a transaction is completed and acknowledged by the database system, it has indeed been permanently recorded and won’t be lost even if a crash ensues shortly thereafter. For example, if we are recording a cash withdrawal by Bob, we do not want any chance that the debit to his account will disappear in a crash just after he walks out the bank door. A transactional database guarantees that all the updates made by a transaction are logged in permanent storage (i.e., on disk) before the transaction is reported complete.

Another important property of transactional databases is closely related to the notion of atomic updates: when multiple transactions are running concurrently, each one should not be able to see the incomplete changes made by others. For example, if one transaction is busy totalling all the branch balances, it would not do for it to include the debit from Alice’s branch but not the credit to Bob’s branch, nor vice versa. So transactions must be all-or-nothing not only in terms of their permanent effect on the database, but also in terms of their visibility as they happen. The updates made so far by an open transaction are invisible to other transactions until the transaction completes, whereupon all the updates become visible simultaneously.

In PostgreSQL, a transaction is set up by surrounding the SQL commands of the transaction with BEGIN and COMMIT commands. So our banking transaction would actually look like:

```
BEGIN;
UPDATE accounts SET balance = balance - 100.00
  WHERE name = 'Alice';
-- etc etc
COMMIT;
```

If, partway through the transaction, we decide we do not want to commit (perhaps we just noticed that Alice’s balance went negative), we can issue the command ROLLBACK instead of COMMIT, and all our updates so far will be canceled.

PostgreSQL actually treats every SQL statement as being executed within a transaction. If you do not issue a BEGIN command, then each individual statement has an implicit BEGIN and (if successful) COMMIT wrapped around it. A group of statements surrounded by BEGIN and COMMIT is sometimes called a transaction block.

**Note:** Some client libraries issue BEGIN and COMMIT commands automatically, so that you might get the effect of transaction blocks without asking. Check the documentation for the interface you are using.

It’s possible to control the statements in a transaction in a more granular fashion through the use of savepoints. Savepoints allow you to selectively discard parts of the transaction, while committing the rest. After defining a savepoint with SAVEPOINT, you can if needed roll back to the savepoint with ROLLBACK TO. All the transaction’s database changes between defining the savepoint and rolling back to it are discarded, but changes earlier than the savepoint are kept.

After rolling back to a savepoint, it continues to be defined, so you can roll back to it several times. Conversely, if you are sure you won’t need to roll back to a particular savepoint again, it can be released, so the system can free some resources. Keep in mind that either releasing or rolling back to a savepoint will automatically release all savepoints that were defined after it.

All this is happening within the transaction block, so none of it is visible to other database sessions. When and if you commit the transaction block, the committed actions become visible as a unit to other sessions, while the rolled-back actions never become visible at all.
Chapter 3. Advanced Features

Remembering the bank database, suppose we debit $100.00 from Alice’s account, and credit Bob’s account, only to find later that we should have credited Wally’s account. We could do it using savepoints like this:

BEGIN;
UPDATE accounts SET balance = balance - 100.00
  WHERE name = 'Alice';
SAVEPOINT my_savepoint;
UPDATE accounts SET balance = balance + 100.00
  WHERE name = 'Bob';
-- oops ... forget that and use Wally’s account
ROLLBACK TO my_savepoint;
UPDATE accounts SET balance = balance + 100.00
  WHERE name = 'Wally';
COMMIT;

This example is, of course, oversimplified, but there’s a lot of control possible in a transaction block through the use of savepoints. Moreover, ROLLBACK TO is the only way to regain control of a transaction block that was put in aborted state by the system due to an error, short of rolling it back completely and starting again.

3.5. Window Functions

A window function performs a calculation across a set of table rows that are somehow related to the current row. This is comparable to the type of calculation that can be done with an aggregate function. But unlike regular aggregate functions, use of a window function does not cause rows to become grouped into a single output row — the rows retain their separate identities. Behind the scenes, the window function is able to access more than just the current row of the query result.

Here is an example that shows how to compare each employee’s salary with the average salary in his or her department:

SELECT depname, empno, salary, avg(salary) OVER (PARTITION BY depname) FROM empsalary;

depname | empno | salary | avg
---------+-------+--------+-----------------------
        |      |        |                       
develop | 11    | 5200   | 5020.0000000000000000
develop | 7     | 4200   | 5020.0000000000000000
develop | 9     | 4500   | 5020.0000000000000000
develop | 8     | 6000   | 5020.0000000000000000
develop | 10    | 5200   | 5020.0000000000000000
personnel | 5     | 3500   | 3700.0000000000000000
personnel | 2     | 3900   | 3700.0000000000000000
sales    | 3     | 4800   | 4866.6666666666666667
sales    | 1     | 5000   | 4866.6666666666666667
sales    | 4     | 4800   | 4866.6666666666666667
(10 rows)

The first three output columns come directly from the table empsalary, and there is one output row for each row in the table. The fourth column represents an average taken across all the table rows that have the same depname value as the current row. (This actually is the same function as the regular avg aggregate function, but the OVER clause causes it to be treated as a window function and computed across an appropriate set of rows.)
A window function call always contains an `OVER` clause directly following the window function’s name and argument(s). This is what syntactically distinguishes it from a regular function or aggregate function. The `OVER` clause determines exactly how the rows of the query are split up for processing by the window function. The `PARTITION BY` list within `OVER` specifies dividing the rows into groups, or partitions, that share the same values of the `PARTITION BY` expression(s). For each row, the window function is computed across the rows that fall into the same partition as the current row.

You can also control the order in which rows are processed by window functions using `ORDER BY` within `OVER`. (The window `ORDER BY` does not even have to match the order in which the rows are output.) Here is an example:

```sql
SELECT depname, empno, salary,
       rank() OVER (PARTITION BY depname ORDER BY salary DESC)
FROM empsalary;
```

<table>
<thead>
<tr>
<th>depname</th>
<th>empno</th>
<th>salary</th>
<th>rank</th>
</tr>
</thead>
<tbody>
<tr>
<td>develop</td>
<td>8</td>
<td>6000</td>
<td>1</td>
</tr>
<tr>
<td>develop</td>
<td>10</td>
<td>5200</td>
<td>2</td>
</tr>
<tr>
<td>develop</td>
<td>11</td>
<td>5200</td>
<td>2</td>
</tr>
<tr>
<td>develop</td>
<td>9</td>
<td>4500</td>
<td>4</td>
</tr>
<tr>
<td>develop</td>
<td>7</td>
<td>4200</td>
<td>5</td>
</tr>
<tr>
<td>personnel</td>
<td>2</td>
<td>3900</td>
<td>1</td>
</tr>
<tr>
<td>personnel</td>
<td>5</td>
<td>3500</td>
<td>2</td>
</tr>
<tr>
<td>sales</td>
<td>1</td>
<td>5000</td>
<td>1</td>
</tr>
<tr>
<td>sales</td>
<td>4</td>
<td>4800</td>
<td>2</td>
</tr>
<tr>
<td>sales</td>
<td>3</td>
<td>4800</td>
<td>2</td>
</tr>
</tbody>
</table>

As shown here, the `rank` function produces a numerical rank within the current row’s partition for each distinct `ORDER BY` value, in the order defined by the `ORDER BY` clause. `rank` needs no explicit parameter, because its behavior is entirely determined by the `OVER` clause.

The rows considered by a window function are those of the “virtual table” produced by the query’s `FROM` clause as filtered by its `WHERE`, `GROUP BY`, and `HAVING` clauses if any. For example, a row removed because it does not meet the `WHERE` condition is not seen by any window function. A query can contain multiple window functions that slice up the data in different ways by means of different `OVER` clauses, but they all act on the same collection of rows defined by this virtual table.

We already saw that `ORDER BY` can be omitted if the ordering of rows is not important. It is also possible to omit `PARTITION BY`, in which case there is just one partition containing all the rows.

There is another important concept associated with window functions: for each row, there is a set of rows within its partition called its `window frame`. Many (but not all) window functions act only on the rows of the window frame, rather than of the whole partition. By default, if `ORDER BY` is supplied then the frame consists of all rows from the start of the partition up through the current row, plus any following rows that are equal to the current row according to the `ORDER BY` clause. When `ORDER BY` is omitted the default frame consists of all rows in the partition. ¹ Here is an example using `sum`:

```sql
SELECT salary, sum(salary) OVER () FROM empsalary;
```

<table>
<thead>
<tr>
<th>salary</th>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>5200</td>
<td>47100</td>
</tr>
<tr>
<td>5000</td>
<td>47100</td>
</tr>
</tbody>
</table>

¹. There are options to define the window frame in other ways, but this tutorial does not cover them. See Section 4.2.8 for details.
Above, since there is no ORDER BY in the OVER clause, the window frame is the same as the partition, which for lack of PARTITION BY is the whole table; in other words each sum is taken over the whole table and so we get the same result for each output row. But if we add an ORDER BY clause, we get very different results:

```sql
SELECT salary, sum(salary) OVER (ORDER BY salary) FROM empsalary;
```

<table>
<thead>
<tr>
<th>salary</th>
<th>sum</th>
</tr>
</thead>
<tbody>
<tr>
<td>3500</td>
<td>3500</td>
</tr>
<tr>
<td>3900</td>
<td>7400</td>
</tr>
<tr>
<td>4200</td>
<td>11600</td>
</tr>
<tr>
<td>4500</td>
<td>16100</td>
</tr>
<tr>
<td>4800</td>
<td>25700</td>
</tr>
<tr>
<td>4800</td>
<td>25700</td>
</tr>
<tr>
<td>5000</td>
<td>30700</td>
</tr>
<tr>
<td>5200</td>
<td>41100</td>
</tr>
<tr>
<td>5200</td>
<td>41100</td>
</tr>
<tr>
<td>6000</td>
<td>47100</td>
</tr>
</tbody>
</table>

Here the sum is taken from the first (lowest) salary up through the current one, including any duplicates of the current one (notice the results for the duplicated salaries).

Window functions are permitted only in the SELECT list and the ORDER BY clause of the query. They are forbidden elsewhere, such as in GROUP BY, HAVING and WHERE clauses. This is because they logically execute after the processing of those clauses. Also, window functions execute after regular aggregate functions. This means it is valid to include an aggregate function call in the arguments of a window function, but not vice versa.

If there is a need to filter or group rows after the window calculations are performed, you can use a sub-select. For example:

```sql
SELECT depname, empno, salary, enroll_date
FROM
(SELECT depname, empno, salary, enroll_date,
    rank() OVER (PARTITION BY depname ORDER BY salary DESC, empno) AS pos
FROM empsalary
) AS ss
WHERE pos < 3;
```

The above query only shows the rows from the inner query having rank less than 3.

When a query involves multiple window functions, it is possible to write out each one with a separate OVER clause, but this is duplicative and error-prone if the same windowing behavior is wanted for several functions. Instead, each windowing behavior can be named in a WINDOW clause and then referenced in OVER. For example:
SELECT sum(salary) OVER w, avg(salary) OVER w
FROM empsalary
WINDOW w AS (PARTITION BY depname ORDER BY salary DESC);

More details about window functions can be found in Section 4.2.8, Section 9.21, Section 7.2.5, and the SELECT reference page.

3.6. Inheritance

Inheritance is a concept from object-oriented databases. It opens up interesting new possibilities of database design.

Let’s create two tables: A table cities and a table capitals. Naturally, capitals are also cities, so you want some way to show the capitals implicitly when you list all cities. If you’re really clever you might invent some scheme like this:

CREATE TABLE capitals {
    name text,
    population real,
    altitude int,    -- (in ft)
    state char(2)
};

CREATE TABLE non_capitals {
    name text,
    population real,
    altitude int    -- (in ft)
};

CREATE VIEW cities AS
    SELECT name, population, altitude FROM capitals
    UNION
    SELECT name, population, altitude FROM non_capitals;

This works OK as far as querying goes, but it gets ugly when you need to update several rows, for one thing.

A better solution is this:

CREATE TABLE cities {
    name text,
    population real,
    altitude int    -- (in ft)
};

CREATE TABLE capitals {
    state char(2)
} INHERITS (cities);

In this case, a row of capitals inherits all columns (name, population, and altitude) from its parent, cities. The type of the column name is text, a native PostgreSQL type for variable length
character strings. State capitals have an extra column, state, that shows their state. In PostgreSQL, a table can inherit from zero or more other tables.

For example, the following query finds the names of all cities, including state capitals, that are located at an altitude over 500 feet:

```sql
SELECT name, altitude
FROM cities
WHERE altitude > 500;
```

which returns:

<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
<tr>
<td>Madison</td>
<td>845</td>
</tr>
</tbody>
</table>

(3 rows)

On the other hand, the following query finds all the cities that are not state capitals and are situated at an altitude over 500 feet:

```sql
SELECT name, altitude
FROM ONLY cities
WHERE altitude > 500;
```

<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
</tbody>
</table>

(2 rows)

Here the ONLY before cities indicates that the query should be run over only the cities table, and not tables below cities in the inheritance hierarchy. Many of the commands that we have already discussed — SELECT, UPDATE, and DELETE — support this ONLY notation.

**Note:** Although inheritance is frequently useful, it has not been integrated with unique constraints or foreign keys, which limits its usefulness. See Section 5.9 for more detail.

### 3.7. Conclusion

PostgreSQL has many features not touched upon in this tutorial introduction, which has been oriented toward newer users of SQL. These features are discussed in more detail in the remainder of this book.

If you feel you need more introductory material, please visit the PostgreSQL web site\(^2\) for links to more resources.

---

\(^2\) [https://www.postgresql.org](https://www.postgresql.org)
II. The SQL Language

This part describes the use of the SQL language in PostgreSQL. We start with describing the general syntax of SQL, then explain how to create the structures to hold data, how to populate the database, and how to query it. The middle part lists the available data types and functions for use in SQL commands. The rest treats several aspects that are important for tuning a database for optimal performance.

The information in this part is arranged so that a novice user can follow it start to end to gain a full understanding of the topics without having to refer forward too many times. The chapters are intended to be self-contained, so that advanced users can read the chapters individually as they choose. The information in this part is presented in a narrative fashion in topical units. Readers looking for a complete description of a particular command should see Part VI.

Readers of this part should know how to connect to a PostgreSQL database and issue SQL commands. Readers that are unfamiliar with these issues are encouraged to read Part I first. SQL commands are typically entered using the PostgreSQL interactive terminal psql, but other programs that have similar functionality can be used as well.
Chapter 4. SQL Syntax

This chapter describes the syntax of SQL. It forms the foundation for understanding the following chapters which will go into detail about how SQL commands are applied to define and modify data.

We also advise users who are already familiar with SQL to read this chapter carefully because it contains several rules and concepts that are implemented inconsistently among SQL databases or that are specific to PostgreSQL.

4.1. Lexical Structure

SQL input consists of a sequence of commands. A command is composed of a sequence of tokens, terminated by a semicolon (";"). The end of the input stream also terminates a command. Which tokens are valid depends on the syntax of the particular command.

A token can be a key word, an identifier, a quoted identifier, a literal (or constant), or a special character symbol. Tokens are normally separated by whitespace (space, tab, newline), but need not be if there is no ambiguity (which is generally only the case if a special character is adjacent to some other token type).

For example, the following is (syntactically) valid SQL input:

```sql
SELECT * FROM MY_TABLE;
UPDATE MY_TABLE SET A = 5;
INSERT INTO MY_TABLE VALUES (3, 'hi there');
```

This is a sequence of three commands, one per line (although this is not required; more than one command can be on a line, and commands can usefully be split across lines).

Additionally, comments can occur in SQL input. They are not tokens, they are effectively equivalent to whitespace.

The SQL syntax is not very consistent regarding what tokens identify commands and which are operands or parameters. The first few tokens are generally the command name, so in the above example we would usually speak of a “SELECT”, an “UPDATE”, and an “INSERT” command. But for instance the UPDATE command always requires a SET token to appear in a certain position, and this particular variation of INSERT also requires a VALUES in order to be complete. The precise syntax rules for each command are described in Part VI.

4.1.1. Identifiers and Key Words

Tokens such as SELECT, UPDATE, or VALUES in the example above are examples of key words, that is, words that have a fixed meaning in the SQL language. The tokens MY_TABLE and A are examples of identifiers. They identify names of tables, columns, or other database objects, depending on the command they are used in. Therefore they are sometimes simply called “names”. Key words and identifiers have the same lexical structure, meaning that one cannot know whether a token is an identifier or a key word without knowing the language. A complete list of key words can be found in Appendix C.

SQL identifiers and key words must begin with a letter (a-z, but also letters with diacritical marks and non-Latin letters) or an underscore (_). Subsequent characters in an identifier or key word can be letters, underscores, digits (0-9), or dollar signs ($). Note that dollar signs are not allowed in identifiers according to the letter of the SQL standard, so their use might render applications less portable. The
Chapter 4. SQL Syntax

SQL standard will not define a key word that contains digits or starts or ends with an underscore, so identifiers of this form are safe against possible conflict with future extensions of the standard.

The system uses no more than NAMEDATALEN−1 bytes of an identifier; longer names can be written in commands, but they will be truncated. By default, NAMEDATALEN is 64 so the maximum identifier length is 63 bytes. If this limit is problematic, it can be raised by changing the NAMEDATALEN constant in src/include/pg_config_manual.h.

Key words and unquoted identifiers are case insensitive. Therefore:

```
UPDATE MY_TABLE SET A = 5;
```

can equivalently be written as:

```
UPDATE my_Table Set a = 5;
```

A convention often used is to write key words in upper case and names in lower case, e.g.:

```
UPDATE my_table SET a = 5;
```

There is a second kind of identifier: the delimited identifier or quoted identifier. It is formed by enclosing an arbitrary sequence of characters in double-quotes (“). A delimited identifier is always an identifier, never a key word. So "select" could be used to refer to a column or table named “select”, whereas an unquoted select would be taken as a key word and would therefore provoke a parse error when used where a table or column name is expected. The example can be written with quoted identifiers like this:

```
UPDATE "my_table" SET "a" = 5;
```

Quoted identifiers can contain any character, except the character with code zero. (To include a double quote, write two double quotes.) This allows constructing table or column names that would otherwise not be possible, such as ones containing spaces or ampersands. The length limitation still applies.

A variant of quoted identifiers allows including escaped Unicode characters identified by their code points. This variant starts with U& (upper or lower case U followed by ampersand) immediately before the opening double quote, without any spaces in between, for example U&"foo". (Note that this creates an ambiguity with the operator &. Use spaces around the operator to avoid this problem.) Inside the quotes, Unicode characters can be specified in escaped form by writing a backslash followed by the four-digit hexadecimal code point number or alternatively a backslash followed by a plus sign followed by a six-digit hexadecimal code point number. For example, the identifier "data" could be written as

```
U&"d\0061t\+000061"
```

The following less trivial example writes the Russian word “слон” (elephant) in Cyrillic letters:

```
U&"\0441\043B\043E\043D"
```

If a different escape character than backslash is desired, it can be specified using the UESCAPE clause after the string, for example:

```
U&"d\0061t\+000061" UESCAPE ‘!’
```
The escape character can be any single character other than a hexadecimal digit, the plus sign, a single quote, a double quote, or a whitespace character. Note that the escape character is written in single quotes, not double quotes.

To include the escape character in the identifier literally, write it twice.

The Unicode escape syntax works only when the server encoding is UTF8. When other server encodings are used, only code points in the ASCII range (up to \007F) can be specified. Both the 4-digit and the 6-digit form can be used to specify UTF-16 surrogate pairs to compose characters with code points larger than U+FFFF, although the availability of the 6-digit form technically makes this unnecessary. (Surrogate pairs are not stored directly, but combined into a single code point that is then encoded in UTF-8.)

Quoting an identifier also makes it case-sensitive, whereas unquoted names are always folded to lower case. For example, the identifiers FOO, foo, and "foo" are considered the same by PostgreSQL, but "Foo" and "FOO" are different from these three and each other. (The folding of unquoted names to lower case in PostgreSQL is incompatible with the SQL standard, which says that unquoted names should be folded to upper case. Thus, foo should be equivalent to "FOO" not "foo" according to the standard. If you want to write portable applications you are advised to always quote a particular name or never quote it.)

### 4.1.2. Constants

There are three kinds of implicitly-typed constants in PostgreSQL: strings, bit strings, and numbers. Constants can also be specified with explicit types, which can enable more accurate representation and more efficient handling by the system. These alternatives are discussed in the following subsections.

#### 4.1.2.1. String Constants

A string constant in SQL is an arbitrary sequence of characters bounded by single quotes ('), for example ‘This is a string’. To include a single-quote character within a string constant, write two adjacent single quotes, e.g., ‘Dianne’s horse’. Note that this is not the same as a double-quote character ("").

Two string constants that are only separated by whitespace *with at least one newline* are concatenated and effectively treated as if the string had been written as one constant. For example:

```
SELECT 'foo'
'bar';
```

is equivalent to:

```
SELECT 'foob';
```

but:

```
SELECT 'foo' 'bar';
```

is not valid syntax. (This slightly bizarre behavior is specified by SQL; PostgreSQL is following the standard.)
4.1.2.2. String Constants with C-style Escapes

PostgreSQL also accepts “escape” string constants, which are an extension to the SQL standard. An escape string constant is specified by writing the letter E (upper or lower case) just before the opening single quote, e.g., E’foo’. (When continuing an escape string constant across lines, write E only before the first opening quote.) Within an escape string, a backslash character (\) begins a C-like backslash escape sequence, in which the combination of backslash and following character(s) represent a special byte value, as shown in Table 4-1.

Table 4-1. Backslash Escape Sequences

<table>
<thead>
<tr>
<th>Backslash Escape Sequence</th>
<th>Interpretation</th>
</tr>
</thead>
<tbody>
<tr>
<td>\b</td>
<td>backspace</td>
</tr>
<tr>
<td>\f</td>
<td>form feed</td>
</tr>
<tr>
<td>\n</td>
<td>newline</td>
</tr>
<tr>
<td>\r</td>
<td>carriage return</td>
</tr>
<tr>
<td>\t</td>
<td>tab</td>
</tr>
<tr>
<td>, \oo, \ooo (o = 0 - 7)</td>
<td>octal byte value</td>
</tr>
<tr>
<td>\hh, \hh (h = 0 - 9, A - F)</td>
<td>hexadecimal byte value</td>
</tr>
<tr>
<td>\xxxxx, \Uxxxxxxxxx (x = 0 - 9, A - F)</td>
<td>16 or 32-bit hexadecimal Unicode character value</td>
</tr>
</tbody>
</table>

Any other character following a backslash is taken literally. Thus, to include a backslash character, write two backslashes (\\). Also, a single quote can be included in an escape string by writing \’, in addition to the normal way of “.

It is your responsibility that the byte sequences you create, especially when using the octal or hexadecimal escapes, compose valid characters in the server character set encoding. When the server encoding is UTF-8, then the Unicode escapes or the alternative Unicode escape syntax, explained in Section 4.1.2.3, should be used instead. (The alternative would be doing the UTF-8 encoding by hand and writing out the bytes, which would be very cumbersome.)

The Unicode escape syntax works fully only when the server encoding is UTF8. When other server encodings are used, only code points in the ASCII range (up to \U007F) can be specified. Both the 4-digit and the 8-digit form can be used to specify UTF-16 surrogate pairs to compose characters with code points larger than U+FFFF, although the availability of the 8-digit form technically makes this unnecessary. (When surrogate pairs are used when the server encoding is UTF8, they are first combined into a single code point that is then encoded in UTF-8.)

Caution

If the configuration parameter standard_conforming_strings is off, then PostgreSQL recognizes backslash escapes in both regular and escape string constants. However, as of PostgreSQL 9.1, the default is on, meaning that backslash escapes are recognized only in escape string constants. This behavior is more standards-compliant, but might break applications which rely on the historical behavior, where backslash escapes were always recognized. As a workaround, you can set this parameter to off, but it is better to migrate away from using backslash escapes. If you need to use a backslash escape to represent a special character, write the string constant with an E.

In addition to standard_conforming_strings, the configuration parameters escape_string_warning and backslash_quote govern treatment of backslashes in string constants.
Chapter 4. SQL Syntax

The character with the code zero cannot be in a string constant.

4.1.2.3. String Constants with Unicode Escapes

PostgreSQL also supports another type of escape syntax for strings that allows specifying arbitrary Unicode characters by code point. A Unicode escape string constant starts with $U$ (upper or lower case letter U followed by ampersand) immediately before the opening quote, without any spaces in between, for example $U'foo'$. (Note that this creates an ambiguity with the operator $&$. Use spaces around the operator to avoid this problem.) Inside the quotes, Unicode characters can be specified in escaped form by writing a backslash followed by the four-digit hexadecimal code point number or alternatively a backslash followed by a plus sign followed by a six-digit hexadecimal code point number. For example, the string ‘data’ could be written as

```
U&'d\0061t\+000061'
```

The following less trivial example writes the Russian word “slon” (elephant) in Cyrillic letters:

```
U&'\0441\043B\043E\043D'
```

If a different escape character than backslash is desired, it can be specified using the `UESCAPE` clause after the string, for example:

```
U&'d!0061t!+000061' UESCAPE '!'  
```

The escape character can be any single character other than a hexadecimal digit, the plus sign, a single quote, a double quote, or a whitespace character.

The Unicode escape syntax works only when the server encoding is UTF8. When other server encodings are used, only code points in the ASCII range (up to \007F) can be specified. Both the 4-digit and the 6-digit form can be used to specify UTF-16 surrogate pairs to compose characters with code points larger than U+FFFF, although the availability of the 6-digit form technically makes this unnecessary. (When surrogate pairs are used when the server encoding is UTF8, they are first combined into a single code point that is then encoded in UTF-8.)

Also, the Unicode escape syntax for string constants only works when the configuration parameter `standard_conforming_strings` is turned on. This is because otherwise this syntax could confuse clients that parse the SQL statements to the point that it could lead to SQL injections and similar security issues. If the parameter is set to off, this syntax will be rejected with an error message.

To include the escape character in the string literally, write it twice.

4.1.2.4. Dollar-quoted String Constants

While the standard syntax for specifying string constants is usually convenient, it can be difficult to understand when the desired string contains many single quotes or backslashes, since each of those must be doubled. To allow more readable queries in such situations, PostgreSQL provides another way, called “dollar quoting”, to write string constants. A dollar-quoted string constant consists of a dollar sign ($), an optional “tag” of zero or more characters, another dollar sign, an arbitrary sequence of characters that makes up the string content, a dollar sign, the same tag that began this dollar quote, and a dollar sign. For example, here are two different ways to specify the string “Dianne’s horse” using dollar quoting:

```
$$Dianne’s horse$$
```
Chapter 4. SQL Syntax

Notice that inside the dollar-quoted string, single quotes can be used without needing to be escaped. Indeed, no characters inside a dollar-quoted string are ever escaped: the string content is always written literally. Backslashes are not special, and neither are dollar signs, unless they are part of a sequence matching the opening tag.

It is possible to nest dollar-quoted string constants by choosing different tags at each nesting level. This is most commonly used in writing function definitions. For example:

```
$function$
BEGIN
  RETURN ($1 ~ $q${[\t\r\n\v\\]}$q$);
END;
$function$
```

Here, the sequence $q{[\t\r\n\v\\]}q$ represents a dollar-quoted literal string {\t\r\n\v\\}, which will be recognized when the function body is executed by PostgreSQL. But since the sequence does not match the outer dollar quoting delimiter $function$, it is just some more characters within the constant so far as the outer string is concerned.

The tag, if any, of a dollar-quoted string follows the same rules as an unquoted identifier, except that it cannot contain a dollar sign. Tags are case sensitive, so $tag$string content$tag$ is correct, but $TAG$string content$tag$ is not.

A dollar-quoted string that follows a keyword or identifier must be separated from it by whitespace; otherwise the dollar quoting delimiter would be taken as part of the preceding identifier.

Dollar quoting is not part of the SQL standard, but it is often a more convenient way to write complicated string literals than the standard-compliant single quote syntax. It is particularly useful when representing string constants inside other constants, as is often needed in procedural function definitions. With single-quote syntax, each backslash in the above example would have to be written as four backslashes, which would be reduced to two backslashes in parsing the original string constant, and then to one when the inner string constant is re-parsed during function execution.

4.1.2.5. Bit-string Constants

Bit-string constants look like regular string constants with a B (upper or lower case) immediately before the opening quote (no intervening whitespace), e.g., B’1001’. The only characters allowed within bit-string constants are 0 and 1.

Alternatively, bit-string constants can be specified in hexadecimal notation, using a leading X (upper or lower case), e.g., X’1FF’. This notation is equivalent to a bit-string constant with four binary digits for each hexadecimal digit.

Both forms of bit-string constant can be continued across lines in the same way as regular string constants. Dollar quoting cannot be used in a bit-string constant.

4.1.2.6. Numeric Constants

Numeric constants are accepted in these general forms:

```
digits
digits.[digits][e[+-]digits]
[digits].[digits][e[+-]digits]
digitse[+-]digits
```
Chapter 4. SQL Syntax

where \textit{digits} is one or more decimal digits (0 through 9). At least one digit must be before or after the decimal point, if one is used. At least one digit must follow the exponent marker (\textit{e}), if one is present. There cannot be any spaces or other characters embedded in the constant. Note that any leading plus or minus sign is not actually considered part of the constant; it is an operator applied to the constant.

These are some examples of valid numeric constants:

\begin{verbatim}
42
3.5
4.
.001
5e2
1.925e-3
\end{verbatim}

A numeric constant that contains neither a decimal point nor an exponent is initially presumed to be type \textit{integer} if its value fits in type \textit{integer} (32 bits); otherwise it is presumed to be type \textit{bigint} if its value fits in type \textit{bigint} (64 bits); otherwise it is taken to be type \textit{numeric}. Constants that contain decimal points and/or exponents are always initially presumed to be type \textit{numeric}.

The initially assigned data type of a numeric constant is just a starting point for the type resolution algorithms. In most cases the constant will be automatically coerced to the most appropriate type depending on context. When necessary, you can force a numeric value to be interpreted as a specific data type by casting it. For example, you can force a numeric value to be treated as type real (float4) by writing:

\begin{verbatim}
REAL '1.23' -- string style
1.23::REAL -- PostgreSQL (historical) style
\end{verbatim}

These are actually just special cases of the general casting notations discussed next.

4.1.2.7. Constants of Other Types

A constant of an \textit{arbitrary} type can be entered using any one of the following notations:

\begin{verbatim}
type 'string'
'string'::type
CAST ( 'string' AS type )
\end{verbatim}

The string constant’s text is passed to the input conversion routine for the type called \textit{type}. The result is a constant of the indicated type. The explicit type cast can be omitted if there is no ambiguity as to the type the constant must be (for example, when it is assigned directly to a table column), in which case it is automatically coerced.

The string constant can be written using either regular SQL notation or dollar-quoting.

It is also possible to specify a type coercion using a function-like syntax:

\begin{verbatim}
typename ( 'string' )
\end{verbatim}

but not all type names can be used in this way; see Section 4.2.9 for details.

The :\textit{,CAST{}}, and function-call syntaxes can also be used to specify run-time type conversions of arbitrary expressions, as discussed in Section 4.2.9. To avoid syntactic ambiguity, the \textit{type 'string'} syntax can only be used to specify the type of a simple literal constant. Another restriction on the \textit{type
Chapter 4. SQL Syntax

The `string` syntax is that it does not work for array types; use `::` or `CAST()` to specify the type of an array constant.

The `CAST()` syntax conforms to SQL. The type `string` syntax is a generalization of the standard: SQL specifies this syntax only for a few data types, but PostgreSQL allows it for all types. The syntax with `::` is historical PostgreSQL usage, as is the function-call syntax.

4.1.3. Operators

An operator name is a sequence of up to `NAMEDATALEN-1` (63 by default) characters from the following list:

```
+ - * / < > = ~ ! @ # % ^ & | ' ?
```

There are a few restrictions on operator names, however:

- `--` and `/*` cannot appear anywhere in an operator name, since they will be taken as the start of a comment.
- A multiple-character operator name cannot end in `+` or `-`, unless the name also contains at least one of these characters:
  ```
  ~ ! @ # % ^ & | ' ?
  ```
  For example, `@-` is an allowed operator name, but `*-` is not. This restriction allows PostgreSQL to parse SQL-compliant queries without requiring spaces between tokens.

When working with non-SQL-standard operator names, you will usually need to separate adjacent operators with spaces to avoid ambiguity. For example, if you have defined a left unary operator named `@`, you cannot write `X*@Y`; you must write `X* @Y` to ensure that PostgreSQL reads it as two operator names not one.

4.1.4. Special Characters

Some characters that are not alphanumeric have a special meaning that is different from being an operator. Details on the usage can be found at the location where the respective syntax element is described. This section only exists to advise the existence and summarize the purposes of these characters.

- A dollar sign (`$`) followed by digits is used to represent a positional parameter in the body of a function definition or a prepared statement. In other contexts the dollar sign can be part of an identifier or a dollar-quoted string constant.
- Parentheses (`()`) have their usual meaning to group expressions and enforce precedence. In some cases parentheses are required as part of the fixed syntax of a particular SQL command.
- Brackets (`[]`) are used to select the elements of an array. See Section 8.15 for more information on arrays.
- Commas (`,`) are used in some syntactical constructs to separate the elements of a list.
• The semicolon (\;) terminates an SQL command. It cannot appear anywhere within a command, except within a string constant or quoted identifier.

• The colon (:\;) is used to select “slices” from arrays. (See Section 8.15.) In certain SQL dialects (such as Embedded SQL), the colon is used to prefix variable names.

• The asterisk (*\;) is used in some contexts to denote all the fields of a table row or composite value. It also has a special meaning when used as the argument of an aggregate function, namely that the aggregate does not require any explicit parameter.

• The period (\.) is used in numeric constants, and to separate schema, table, and column names.

4.1.5. Comments

A comment is a sequence of characters beginning with double dashes and extending to the end of the line, e.g.:

-- This is a standard SQL comment

Alternatively, C-style block comments can be used:

/* multiline comment
 * with nesting: /* nested block comment */
 */

where the comment begins with /* and extends to the matching occurrence of */. These block comments nest, as specified in the SQL standard but unlike C, so that one can comment out larger blocks of code that might contain existing block comments.

A comment is removed from the input stream before further syntax analysis and is effectively replaced by whitespace.

4.1.6. Operator Precedence

Table 4-2 shows the precedence and associativity of the operators in PostgreSQL. Most operators have the same precedence and are left-associative. The precedence and associativity of the operators is hard-wired into the parser.

You will sometimes need to add parentheses when using combinations of binary and unary operators. For instance:

SELECT 5 ! - 6;

will be parsed as:

SELECT 5 ! (- 6);

because the parser has no idea — until it is too late — that ! is defined as a postfix operator, not an infix one. To get the desired behavior in this case, you must write:

SELECT (5 !) - 6;

This is the price one pays for extensibility.
Table 4-2. Operator Precedence (highest to lowest)

<table>
<thead>
<tr>
<th>Operator/Element</th>
<th>Associativity</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>.</td>
<td>left</td>
<td>table/column name separator</td>
</tr>
<tr>
<td>::</td>
<td>left</td>
<td>PostgreSQL-style typecast</td>
</tr>
<tr>
<td>[ ]</td>
<td>left</td>
<td>array element selection</td>
</tr>
<tr>
<td>+ -</td>
<td>right</td>
<td>unary plus, unary minus</td>
</tr>
<tr>
<td>^</td>
<td>left</td>
<td>exponentiation</td>
</tr>
<tr>
<td>* / %</td>
<td>left</td>
<td>multiplication, division, modulo</td>
</tr>
<tr>
<td>+ -</td>
<td>left</td>
<td>addition, subtraction</td>
</tr>
<tr>
<td>(any other operator)</td>
<td>left</td>
<td>all other native and user-defined operators</td>
</tr>
<tr>
<td>BETWEEN IN LIKE ILIKE SIMILAR</td>
<td></td>
<td>range containment, set membership, string matching</td>
</tr>
<tr>
<td>&lt; &gt; &lt;= &gt;= &lt; &gt; &lt; &gt;</td>
<td></td>
<td>comparison operators</td>
</tr>
<tr>
<td>IS ISNULL NOTNULL</td>
<td></td>
<td>IS TRUE, IS FALSE, IS NULL, IS DISTINCT FROM, etc</td>
</tr>
<tr>
<td>NOT</td>
<td>right</td>
<td>logical negation</td>
</tr>
<tr>
<td>AND</td>
<td>left</td>
<td>logical conjunction</td>
</tr>
<tr>
<td>OR</td>
<td>left</td>
<td>logical disjunction</td>
</tr>
</tbody>
</table>

Note that the operator precedence rules also apply to user-defined operators that have the same names as the built-in operators mentioned above. For example, if you define a “+” operator for some custom data type it will have the same precedence as the built-in “+” operator, no matter what yours does.

When a schema-qualified operator name is used in the OPERATOR syntax, as for example in:

```
SELECT 3 OPERATOR(pg_catalog.+) 4;
```

the OPERATOR construct is taken to have the default precedence shown in Table 4-2 for “any other operator”. This is true no matter which specific operator appears inside OPERATOR().

**Note:** PostgreSQL versions before 9.5 used slightly different operator precedence rules. In particular, <= = >= and <> used to be treated as generic operators; IS tests used to have higher priority; and NOT BETWEEN and related constructs acted inconsistently, being taken in some cases as having the precedence of NOT rather than BETWEEN. These rules were changed for better compliance with the SQL standard and to reduce confusion from inconsistent treatment of logically equivalent constructs. In most cases, these changes will result in no behavioral change, or perhaps in “no such operator” failures which can be resolved by adding parentheses. However there are corner cases in which a query might change behavior without any parsing error being reported. If you are concerned about whether these changes have silently broken something, you can test your application with the configuration parameter operator_precedence_warning turned on to see if any warnings are logged.
4.2. Value Expressions

Value expressions are used in a variety of contexts, such as in the target list of the `SELECT` command, as new column values in `INSERT` or `UPDATE`, or in search conditions in a number of commands. The result of a value expression is sometimes called a `scalar`, to distinguish it from the result of a table expression (which is a table). Value expressions are therefore also called `scalar expressions` (or even simply `expressions`). The expression syntax allows the calculation of values from primitive parts using arithmetic, logical, set, and other operations.

A value expression is one of the following:

- A constant or literal value
- A column reference
- A positional parameter reference, in the body of a function definition or prepared statement
- A subscripted expression
- A field selection expression
- An operator invocation
- A function call
- An aggregate expression
- A window function call
- A type cast
- A collation expression
- A scalar subquery
- An array constructor
- A row constructor
- Another value expression in parentheses (used to group subexpressions and override precedence)

In addition to this list, there are a number of constructs that can be classified as an expression but do not follow any general syntax rules. These generally have the semantics of a function or operator and are explained in the appropriate location in Chapter 9. An example is the `IS NULL` clause.

We have already discussed constants in Section 4.1.2. The following sections discuss the remaining options.

4.2.1. Column References

A column can be referenced in the form:

```
correlation.columnname
```

`correlation` is the name of a table (possibly qualified with a schema name), or an alias for a table defined by means of a `FROM` clause. The correlation name and separating dot can be omitted if the column name is unique across all the tables being used in the current query. (See also Chapter 7.)
4.2.2. Positional Parameters

A positional parameter reference is used to indicate a value that is supplied externally to an SQL statement. Parameters are used in SQL function definitions and in prepared queries. Some client libraries also support specifying data values separately from the SQL command string, in which case parameters are used to refer to the out-of-line data values. The form of a parameter reference is:

$number$

For example, consider the definition of a function, `dept`, as:

```
CREATE FUNCTION dept(text) RETURNS dept
    AS $$ SELECT * FROM dept WHERE name = $1 $$
    LANGUAGE SQL;
```

Here the `$1` references the value of the first function argument whenever the function is invoked.

4.2.3. Subscripts

If an expression yields a value of an array type, then a specific element of the array value can be extracted by writing

`expression[subscript]`

or multiple adjacent elements (an “array slice”) can be extracted by writing

`expression[lower_subscript:upper_subscript]`

(Here, the brackets `[ ]` are meant to appear literally.) Each subscript is itself an expression, which must yield an integer value.

In general the array `expression` must be parenthesized, but the parentheses can be omitted when the expression to be subscripted is just a column reference or positional parameter. Also, multiple subscripts can be concatenated when the original array is multidimensional. For example:

```
mytable.arraycolumn[4]
mytable.two_d_column[17][34]
$1[10:42]
(arrayfunction(a,b))[42]
```

The parentheses in the last example are required. See Section 8.15 for more about arrays.

4.2.4. Field Selection

If an expression yields a value of a composite type (row type), then a specific field of the row can be extracted by writing

`expression.fieldname`

In general the row `expression` must be parenthesized, but the parentheses can be omitted when the expression to be selected from is just a table reference or positional parameter. For example:
mytable.mycolumn
$1.somecolumn
(rowfunction(a,b)).col3

(Thus, a qualified column reference is actually just a special case of the field selection syntax.) An important special case is extracting a field from a table column that is of a composite type:

(compositecol).somefield
(mytable.compositecol).somefield

The parentheses are required here to show that compositecol is a column name not a table name, or that mytable is a table name not a schema name in the second case.

You can ask for all fields of a composite value by writing .*

(compositecol).*

This notation behaves differently depending on context; see Section 8.16.5 for details.

4.2.5. Operator Invocations

There are three possible syntaxes for an operator invocation:

expression operator expression (binary infix operator)
operator expression (unary prefix operator)
expression operator (unary postfix operator)

where the operator token follows the syntax rules of Section 4.1.3, or is one of the key words AND, OR, and NOT, or is a qualified operator name in the form:

OPERATOR(schema.operatorname)

Which particular operators exist and whether they are unary or binary depends on what operators have been defined by the system or the user. Chapter 9 describes the built-in operators.

4.2.6. Function Calls

The syntax for a function call is the name of a function (possibly qualified with a schema name), followed by its argument list enclosed in parentheses:

function_name ([expression [, expression ... ]])

For example, the following computes the square root of 2:

sqrt(2)

The list of built-in functions is in Chapter 9. Other functions can be added by the user.

When issuing queries in a database where some users mistrust other users, observe security precautions from Section 10.3 when writing function calls.

The arguments can optionally have names attached. See Section 4.3 for details.
Note: A function that takes a single argument of composite type can optionally be called using field-selection syntax, and conversely field selection can be written in functional style. That is, the notations `col(table)` and `table.col` are interchangeable. This behavior is not SQL-standard but is provided in PostgreSQL because it allows use of functions to emulate “computed fields”. For more information see Section 8.16.5.

4.2.7. Aggregate Expressions

An aggregate expression represents the application of an aggregate function across the rows selected by a query. An aggregate function reduces multiple inputs to a single output value, such as the sum or average of the inputs. The syntax of an aggregate expression is one of the following:

```
aggregate_name (expression [ , ... ] [ order_by_clause ] [ FILTER ( WHERE filter_clause ) ] )
aggregate_name (ALL expression [ , ... ] [ order_by_clause ] [ FILTER ( WHERE filter_clause ) ] )
aggregate_name (DISTINCT expression [ , ... ] [ order_by_clause ] [ FILTER ( WHERE filter_clause ) ] )
aggregate_name ( * ) [ FILTER ( WHERE filter_clause ) ]
aggregate_name ( [ expression [ , ... ] ] ) WITHIN GROUP ( order_by_clause ) [ FILTER ( WHERE filter_clause ) ]
```

where `aggregate_name` is a previously defined aggregate (possibly qualified with a schema name) and `expression` is any value expression that does not itself contain an aggregate expression or a window function call. The optional `order_by_clause` and `filter_clause` are described below.

The first form of aggregate expression invokes the aggregate once for each input row. The second form is the same as the first, since `ALL` is the default. The third form invokes the aggregate once for each distinct value of the expression (or distinct set of values, for multiple expressions) found in the input rows. The fourth form invokes the aggregate once for each input row; since no particular input value is specified, it is generally only useful for the `count(*)` aggregate function. The last form is used with ordered-set aggregate functions, which are described below.

Most aggregate functions ignore null inputs, so that rows in which one or more of the expression(s) yield null are discarded. This can be assumed to be true, unless otherwise specified, for all built-in aggregates.

For example, `count(*)` yields the total number of input rows; `count(f1)` yields the number of input rows in which `f1` is non-null, since `count` ignores nulls; and `count(distinct f1)` yields the number of distinct non-null values of `f1`.

Ordinarily, the input rows are fed to the aggregate function in an unspecified order. In many cases this does not matter; for example, `min` produces the same result no matter what order it receives the inputs in. However, some aggregate functions (such as `array_agg` and `string_agg`) produce results that depend on the ordering of the input rows. When using such an aggregate, the optional `order_by_clause` can be used to specify the desired ordering. The `order_by_clause` has the same syntax as for a query-level `ORDER BY` clause, as described in Section 7.5, except that its expressions are always just expressions and cannot be output-column names or numbers. For example:

```
SELECT array_agg(a ORDER BY b DESC) FROM table;
```

When dealing with multiple-argument aggregate functions, note that the `ORDER BY` clause goes after all the aggregate arguments. For example, write this:

```
SELECT string_agg(a, ',') ORDER BY a FROM table;
```
not this:

```
SELECT string_agg(a ORDER BY a, ',') FROM table; -- incorrect
```

The latter is syntactically valid, but it represents a call of a single-argument aggregate function with two `ORDER BY` keys (the second one being rather useless since it’s a constant).

If `DISTINCT` is specified in addition to an `ORDER BY` clause, then all the `ORDER BY` expressions must match regular arguments of the aggregate; that is, you cannot sort on an expression that is not included in the `DISTINCT` list.

**Note:** The ability to specify both `DISTINCT` and `ORDER BY` in an aggregate function is a PostgreSQL extension.

Placing `ORDER BY` within the aggregate’s regular argument list, as described so far, is used when ordering the input rows for a “normal” aggregate for which ordering is optional. There is a subclass of aggregate functions called **ordered-set aggregates** for which an `ORDER BY` clause is required, usually because the aggregate’s computation is only sensible in terms of a specific ordering of its input rows. Typical examples of ordered-set aggregates include rank and percentile calculations. For an ordered-set aggregate, the `ORDER BY` clause is written inside `WITHIN GROUP (...)`, as shown in the final syntax alternative above. The expressions in the `ORDER BY` clause are evaluated once per input row just like normal aggregate arguments, sorted as per the `ORDER BY` clause’s requirements, and fed to the aggregate function as input arguments. (This is unlike the case for a non-`WITHIN GROUP` `ORDER BY` clause, which is not treated as argument(s) to the aggregate function.) The argument expressions preceding `WITHIN GROUP`, if any, are called **direct arguments** to distinguish them from the **aggregated arguments** listed in the `ORDER BY` clause. Unlike normal aggregate arguments, direct arguments are evaluated only once per aggregate call, not once per input row. This means that they can contain variables only if those variables are grouped by `GROUP BY`; this restriction is the same as if the direct arguments were not inside an aggregate expression at all. Direct arguments are typically used for things like percentile fractions, which only make sense as a single value per aggregation calculation. The direct argument list can be empty; in this case, write just `()` not `(*)`. (PostgreSQL will actually accept either spelling, but only the first way conforms to the SQL standard.)

An example of an ordered-set aggregate call is:

```
SELECT percentile_cont(0.5) WITHIN GROUP (ORDER BY income) FROM households;
```

```
percentile_cont
---------------
      50489
```

which obtains the 50th percentile, or median, value of the `income` column from table `households`. Here, `0.5` is a direct argument; it would make no sense for the percentile fraction to be a value varying across rows.

If `FILTER` is specified, then only the input rows for which the `FILTER` clause evaluates to true are fed to the aggregate function; other rows are discarded. For example:

```
SELECT count(*) AS unfiltered,
       count(*) FILTER (WHERE i < 5) AS filtered
FROM generate_series(1,10) AS s(i);
```

```
unfiltered | filtered
------------+----------
       10 |         4

(1 row)
```
The predefined aggregate functions are described in Section 9.20. Other aggregate functions can be added by the user.

An aggregate expression can only appear in the result list or HAVING clause of a SELECT command. It is forbidden in other clauses, such as WHERE, because those clauses are logically evaluated before the results of aggregates are formed.

When an aggregate expression appears in a subquery (see Section 4.2.11 and Section 9.22), the aggregate is normally evaluated over the rows of the subquery. But an exception occurs if the aggregate’s arguments (and filter_clause if any) contain only outer-level variables: the aggregate then belongs to the nearest such outer level, and is evaluated over the rows of that query. The aggregate expression as a whole is then an outer reference for the subquery it appears in, and acts as a constant over any one evaluation of that subquery. The restriction about appearing only in the result list or HAVING clause applies with respect to the query level that the aggregate belongs to.

### 4.2.8. Window Function Calls

A window function call represents the application of an aggregate-like function over some portion of the rows selected by a query. Unlike regular aggregate function calls, this is not tied to grouping of the selected rows into a single output row — each row remains separate in the query output. However the window function is able to scan all the rows that would be part of the current row’s group according to the grouping specification (PARTITION BY list) of the window function call. The syntax of a window function call is one of the following:

```
function_name ([expression [, expression ... ]]) [ FILTER ( WHERE filter_clause ) ] OVER window_name
```

```
function_name ( * ) [ FILTER ( WHERE filter_clause ) ] OVER window_name
```

where `window_definition` has the syntax

```
[ existing_window_name ]
[ PARTITION BY expression [, ...] ]
[ ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS { FIRST | LAST } ] [, ...] ]
[ frame_clause ]
```

and the optional `frame_clause` can be one of

```
{ RANGE | ROWS } frame_start
{ RANGE | ROWS } BETWEEN frame_start AND frame_end
```

where `frame_start` and `frame_end` can be one of

- UNBOUNDED PRECEDING
- value PRECEDING
- CURRENT ROW
- value FOLLOWING
- UNBOUNDED FOLLOWING

Here, `expression` represents any value expression that does not itself contain window function calls. `window_name` is a reference to a named window specification defined in the query’s WINDOW clause. Alternatively, a full `window_definition` can be given within parentheses, using the same syntax.
as for defining a named window in the WINDOW clause; see the SELECT reference page for details. It’s worth pointing out that OVER wname is not exactly equivalent to OVER (wname); the latter implies copying and modifying the window definition, and will be rejected if the referenced window specification includes a frame clause.

The PARTITION BY option groups the rows of the query into partitions, which are processed separately by the window function. PARTITION BY works similarly to a query-level GROUP BY clause, except that its expressions are always just expressions and cannot be output-column names or numbers. Without PARTITION BY, all rows produced by the query are treated as a single partition. The ORDER BY option determines the order in which the rows of a partition are processed by the window function. It works similarly to a query-level ORDER BY clause, but likewise cannot use output-column names or numbers. Without ORDER BY, rows are processed in an unspecified order.

The frame_clause specifies the set of rows constituting the window frame, which is a subset of the current partition, for those window functions that act on the frame instead of the whole partition. The frame can be specified in either RANGE or ROWS mode; in either case, it runs from the frame_start to the frame_end. If frame_end is omitted, it defaults to CURRENT ROW.

A frame_start of UNBOUNDED PRECEDING means that the frame starts with the first row of the partition, and similarly a frame_end of UNBOUNDED FOLLOWING means that the frame ends with the last row of the partition.

In RANGE mode, a frame_start of CURRENT ROW means the frame starts with the current row’s first peer row (a row that ORDER BY considers equivalent to the current row), while a frame_end of CURRENT ROW means the frame ends with the last equivalent ORDER BY peer. In ROWS mode, CURRENT ROW simply means the current row.

The value PRECEDING and value FOLLOWING cases are currently only allowed in ROWS mode. They indicate that the frame starts or ends the specified number of rows before or after the current row. value must be an integer expression not containing any variables, aggregate functions, or window functions. The value must not be null or negative; but it can be zero, which just selects the current row.

The default framing option is RANGE UNBOUNDED PRECEDING, which is the same as RANGE BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW. With ORDER BY, this sets the frame to be all rows from the partition start up through the current row’s last ORDER BY peer. Without ORDER BY, all rows of the partition are included in the window frame, since all rows become peers of the current row.

Restrictions are that frame_start cannot be UNBOUNDED FOLLOWING, frame_end cannot be UNBOUNDED PRECEDING, and the frame_end choice cannot appear earlier in the above list than the frame_start choice — for example RANGE BETWEEN CURRENT ROW AND value PRECEDING is not allowed.

If FILTER is specified, then only the input rows for which the filter_clause evaluates to true are fed to the window function; other rows are discarded. Only window functions that are aggregates accept a FILTER clause.

The built-in window functions are described in Table 9-56. Other window functions can be added by the user. Also, any built-in or user-defined normal aggregate function can be used as a window function. Ordered-set aggregates presently cannot be used as window functions, however.

The syntaxes using * are used for calling parameter-less aggregate functions as window functions, for example count(*) OVER (PARTITION BY x ORDER BY y). The asterisk (*) is customarily not used for non-aggregate window functions. Aggregate window functions, unlike normal aggregate functions, do not allow DISTINCT or ORDER BY to be used within the function argument list.

Window function calls are permitted only in the SELECT list and the ORDER BY clause of the query.
More information about window functions can be found in Section 3.5, Section 9.21, and Section 7.2.5.

### 4.2.9. Type Casts

A type cast specifies a conversion from one data type to another. PostgreSQL accepts two equivalent syntaxes for type casts:

\[
\text{CAST (} \text{expression AS type } \text{)} \quad \text{expression::type}
\]

The \text{CAST} syntax conforms to SQL; the syntax with :: is historical PostgreSQL usage.

When a cast is applied to a value expression of a known type, it represents a run-time type conversion. The cast will succeed only if a suitable type conversion operation has been defined. Notice that this is subtly different from the use of casts with constants, as shown in Section 4.1.2.7. A cast applied to an unadorned string literal represents the initial assignment of a type to a literal constant value, and so it will succeed for any type (if the contents of the string literal are acceptable input syntax for the data type).

An explicit type cast can usually be omitted if there is no ambiguity as to the type that a value expression must produce (for example, when it is assigned to a table column); the system will automatically apply a type cast in such cases. However, automatic casting is only done for casts that are marked “OK to apply implicitly” in the system catalogs. Other casts must be invoked with explicit casting syntax. This restriction is intended to prevent surprising conversions from being applied silently.

It is also possible to specify a type cast using a function-like syntax:

\[
\text{typename (} \text{expression } \text{)}
\]

However, this only works for types whose names are also valid as function names. For example, \text{double precision} cannot be used this way, but the equivalent \text{float8} can. Also, the names \text{interval}, \text{time}, and \text{timestamp} can only be used in this fashion if they are double-quoted, because of syntactic conflicts. Therefore, the use of the function-like cast syntax leads to inconsistencies and should probably be avoided.

**Note:** The function-like syntax is in fact just a function call. When one of the two standard cast syntaxes is used to do a run-time conversion, it will internally invoke a registered function to perform the conversion. By convention, these conversion functions have the same name as their output type, and thus the “function-like syntax” is nothing more than a direct invocation of the underlying conversion function. Obviously, this is not something that a portable application should rely on. For further details see \text{CREATE CAST}.

### 4.2.10. Collation Expressions

The \text{COLLATE} clause overrides the collation of an expression. It is appended to the expression it applies to:

\[
\text{expr COLLATE collation}
\]

where \text{collation} is a possibly schema-qualified identifier. The \text{COLLATE} clause binds tighter than operators; parentheses can be used when necessary.
If no collation is explicitly specified, the database system either derives a collation from the columns involved in the expression, or it defaults to the default collation of the database if no column is involved in the expression.

The two common uses of the \texttt{COLLATE} clause are overriding the sort order in an \texttt{ORDER BY} clause, for example:

\begin{verbatim}
SELECT a, b, c FROM tbl WHERE ... ORDER BY a COLLATE "C";
\end{verbatim}

and overriding the collation of a function or operator call that has locale-sensitive results, for example:

\begin{verbatim}
SELECT * FROM tbl WHERE a > 'foo' COLLATE "C";
\end{verbatim}

Note that in the latter case the \texttt{COLLATE} clause is attached to an input argument of the operator we wish to affect. It doesn’t matter which argument of the operator or function call the \texttt{COLLATE} clause is attached to, because the collation that is applied by the operator or function is derived by considering all arguments, and an explicit \texttt{COLLATE} clause will override the collations of all other arguments. (Attaching non-matching \texttt{COLLATE} clauses to more than one argument, however, is an error. For more details see Section 23.2.) Thus, this gives the same result as the previous example:

\begin{verbatim}
SELECT * FROM tbl WHERE a COLLATE "C" > 'foo';
\end{verbatim}

But this is an error:

\begin{verbatim}
SELECT * FROM tbl WHERE (a > 'foo') COLLATE "C";
\end{verbatim}

because it attempts to apply a collation to the result of the \texttt{>} operator, which is of the non-collatable data type \texttt{boolean}.

### 4.2.11. Scalar Subqueries

A scalar subquery is an ordinary \texttt{SELECT} query in parentheses that returns exactly one row with one column. (See Chapter 7 for information about writing queries.) The \texttt{SELECT} query is executed and the single returned value is used in the surrounding value expression. It is an error to use a query that returns more than one row or more than one column as a scalar subquery. (But if, during a particular execution, the subquery returns no rows, there is no error; the scalar result is taken to be null.) The subquery can refer to variables from the surrounding query, which will act as constants during any one evaluation of the subquery. See also Section 9.22 for other expressions involving subqueries.

For example, the following finds the largest city population in each state:

\begin{verbatim}
SELECT name, (SELECT max(pop) FROM cities WHERE cities.state = states.name) FROM states;
\end{verbatim}

### 4.2.12. Array Constructors

An array constructor is an expression that builds an array value using values for its member elements. A simple array constructor consists of the key word \texttt{ARRAY}, a left square bracket \texttt{[}, a list of expressions (separated by commas) for the array element values, and finally a right square bracket \texttt{]}. For example:

\begin{verbatim}
SELECT ARRAY[1,2,3+4];
\end{verbatim}
By default, the array element type is the common type of the member expressions, determined using the same rules as for UNION or CASE constructs (see Section 10.5). You can override this by explicitly casting the array constructor to the desired type, for example:

```sql
SELECT ARRAY[1,2,22.7]::integer[];
```

This has the same effect as casting each expression to the array element type individually. For more on casting, see Section 4.2.9.

Multidimensional array values can be built by nesting array constructors. In the inner constructors, the key word ARRAY can be omitted. For example, these produce the same result:

```sql
SELECT ARRAY[ARRAY[1,2], ARRAY[3,4]];
```

Since multidimensional arrays must be rectangular, inner constructors at the same level must produce sub-arrays of identical dimensions. Any cast applied to the outer ARRAY constructor propagates automatically to all the inner constructors.

Multidimensional array constructor elements can be anything yielding an array of the proper kind, not only a sub-ARRAY construct. For example:

```sql
CREATE TABLE arr(f1 int[], f2 int[]);
```

```sql
SELECT ARRAY[f1, f2, '{{9,10},{11,12}}'::int[]] FROM arr;
```

You can construct an empty array, but since it’s impossible to have an array with no type, you must explicitly cast your empty array to the desired type. For example:

```sql
SELECT ARRAY[]::integer[];
```
It is also possible to construct an array from the results of a subquery. In this form, the array constructor is written with the key word `ARRAY` followed by a parenthesized (not bracketed) subquery. For example:

```sql
SELECT ARRAY(SELECT oid FROM pg_proc WHERE proname LIKE 'bytea%');
```

```
array
```

```sql
SELECT ARRAY(SELECT ARRAY[i, i*2] FROM generate_series(1,5) AS a(i));
```

```
array
{{1,2},{2,4},{3,6},{4,8},{5,10}}
```

The subquery must return a single column. If the subquery's output column is of a non-array type, the resulting one-dimensional array will have an element for each row in the subquery result, with an element type matching that of the subquery's output column. If the subquery's output column is of an array type, the result will be an array of the same type but one higher dimension; in this case all the subquery rows must yield arrays of identical dimensionality, else the result would not be rectangular.

The subscripts of an array value built with `ARRAY` always begin with one. For more information about arrays, see Section 8.15.

### 4.2.13. Row Constructors

A row constructor is an expression that builds a row value (also called a composite value) using values for its member fields. A row constructor consists of the key word `ROW`, a left parenthesis, zero or more expressions (separated by commas) for the row field values, and finally a right parenthesis. For example:

```sql
SELECT ROW(1,2.5,'this is a test');
```

The key word `ROW` is optional when there is more than one expression in the list.

A row constructor can include the syntax `rowvalue.*,` which will be expanded to a list of the elements of the row value, just as occurs when the `.*` syntax is used at the top level of a `SELECT` list (see Section 8.16.5). For example, if table `t` has columns `f1` and `f2`, these are the same:

```sql
SELECT ROW(t.*, 42) FROM t;
SELECT ROW(t.f1, t.f2, 42) FROM t;
```

**Note:** Before PostgreSQL 8.2, the `.*` syntax was not expanded in row constructors, so that writing `ROW(t.*, 42)` created a two-field row whose first field was another row value. The new behavior is usually more useful. If you need the old behavior of nested row values, write the inner row value without `.*`, for instance `ROW(t, 42).`
By default, the value created by a \texttt{ROW} expression is of an anonymous record type. If necessary, it can be cast to a named composite type — either the row type of a table, or a composite type created with \texttt{CREATE TYPE AS}. An explicit cast might be needed to avoid ambiguity. For example:

\begin{verbatim}
CREATE TABLE mytable(f1 int, f2 float, f3 text);

CREATE FUNCTION getf1(mytable) RETURNS int AS 'SELECT $1.f1' LANGUAGE SQL;

-- No cast needed since only one getf1() exists
SELECT getf1(ROW(1,2.5,'this is a test'));
getf1
-------
 1
(1 row)

CREATE TYPE myrowtype AS (f1 int, f2 text, f3 numeric);

CREATE FUNCTION getf1(myrowtype) RETURNS int AS 'SELECT $1.f1' LANGUAGE SQL;

-- Now we need a cast to indicate which function to call:
SELECT getf1(ROW(1,2.5,'this is a test'));
ERROR: function getf1(record) is not unique

SELECT getf1(ROW(1,2.5,'this is a test')::mytable);
getf1
-------
 1
(1 row)

SELECT getf1(CAST(ROW(11,'this is a test',2.5) AS myrowtype));
getf1
-------
 11
(1 row)
\end{verbatim}

Row constructors can be used to build composite values to be stored in a composite-type table column, or to be passed to a function that accepts a composite parameter. Also, it is possible to compare two row values or test a row with \texttt{IS NULL} or \texttt{IS NOT NULL}, for example:

\begin{verbatim}
SELECT ROW(1,2.5,'this is a test') = ROW(1, 3, 'not the same');
SELECT ROW(table.*) IS NULL FROM table; -- detect all-null rows
\end{verbatim}

For more detail see Section 9.23. Row constructors can also be used in connection with subqueries, as discussed in Section 9.22.

### 4.2.14. Expression Evaluation Rules

The order of evaluation of subexpressions is not defined. In particular, the inputs of an operator or function are not necessarily evaluated left-to-right or in any other fixed order.

Furthermore, if the result of an expression can be determined by evaluating only some parts of it, then other subexpressions might not be evaluated at all. For instance, if one wrote:
Chapter 4. SQL Syntax

SELECT true OR somefunc();

then somefunc() would (probably) not be called at all. The same would be the case if one wrote:

SELECT somefunc() OR true;

Note that this is not the same as the left-to-right “short-circuiting” of Boolean operators that is found in some programming languages.

As a consequence, it is unwise to use functions with side effects as part of complex expressions. It is particularly dangerous to rely on side effects or evaluation order in WHERE and HAVING clauses, since those clauses are extensively reprocessed as part of developing an execution plan. Boolean expressions (AND/OR/NOT combinations) in those clauses can be reorganized in any manner allowed by the laws of Boolean algebra.

When it is essential to force evaluation order, a CASE construct (see Section 9.17) can be used. For example, this is an untrustworthy way of trying to avoid division by zero in a WHERE clause:

SELECT ... WHERE x > 0 AND y/x > 1.5;

But this is safe:

SELECT ... WHERE CASE WHEN x > 0 THEN y/x > 1.5 ELSE false END;

A CASE construct used in this fashion will defeat optimization attempts, so it should only be done when necessary. (In this particular example, it would be better to sidestep the problem by writing \( y > 1.5 \times x \) instead.)

CASE is not a cure-all for such issues, however. One limitation of the technique illustrated above is that it does not prevent early evaluation of constant subexpressions. As described in Section 36.6, functions and operators marked IMMUTABLE can be evaluated when the query is planned rather than when it is executed. Thus for example

SELECT CASE WHEN x > 0 THEN x ELSE 1/0 END FROM tab;

is likely to result in a division-by-zero failure due to the planner trying to simplify the constant subexpression, even if every row in the table has \( x > 0 \) so that the ELSE arm would never be entered at run time.

While that particular example might seem silly, related cases that don’t obviously involve constants can occur in queries executed within functions, since the values of function arguments and local variables can be inserted into queries as constants for planning purposes. Within PL/pgSQL functions, for example, using an IF-THEN-ELSE statement to protect a risky computation is much safer than just nesting it in a CASE expression.

Another limitation of the same kind is that a CASE cannot prevent evaluation of an aggregate expression contained within it, because aggregate expressions are computed before other expressions in a SELECT list or HAVING clause are considered. For example, the following query can cause a division-by-zero error despite seemingly having protected against it:

SELECT CASE WHEN min(employees) > 0 THEN avg(expenses / employees) END FROM departments;

The min() and avg() aggregates are computed concurrently over all the input rows, so if any row has employees equal to zero, the division-by-zero error will occur before there is any opportunity
to test the result of \texttt{min()}. Instead, use a \texttt{WHERE} or \texttt{FILTER} clause to prevent problematic input rows from reaching an aggregate function in the first place.

4.3. Calling Functions

PostgreSQL allows functions that have named parameters to be called using either \textit{positional} or \textit{named} notation. Named notation is especially useful for functions that have a large number of parameters, since it makes the associations between parameters and actual arguments more explicit and reliable. In positional notation, a function call is written with its argument values in the same order as they are defined in the function declaration. In named notation, the arguments are matched to the function parameters by name and can be written in any order. For each notation, also consider the effect of function argument types, documented in Section 10.3.

In either notation, parameters that have default values given in the function declaration need not be written in the call at all. But this is particularly useful in named notation, since any combination of parameters can be omitted; while in positional notation parameters can only be omitted from right to left.

PostgreSQL also supports \textit{mixed} notation, which combines positional and named notation. In this case, positional parameters are written first and named parameters appear after them. The following examples will illustrate the usage of all three notations, using the following function definition:

```sql
CREATE FUNCTION concat_lower_or_upper(a text, b text, uppercase boolean DEFAULT false) RETURNS text
AS $$
SELECT CASE
  WHEN $3 THEN UPPER($1 || ' ' || $2)
  ELSE LOWER($1 || ' ' || $2)
END;
$$
LANGUAGE SQL IMMUTABLE STRICT;
```

Function \texttt{concat_lower_or_upper} has two mandatory parameters, \texttt{a} and \texttt{b}. Additionally there is one optional parameter \texttt{uppercase} which defaults to \texttt{false}. The \texttt{a} and \texttt{b} inputs will be concatenated, and forced to either upper or lower case depending on the \texttt{uppercase} parameter. The remaining details of this function definition are not important here (see Chapter 36 for more information).

4.3.1. Using Positional Notation

Positional notation is the traditional mechanism for passing arguments to functions in PostgreSQL. An example is:

```sql
SELECT concat_lower_or_upper('Hello', 'World', true);
```

```
HELLO WORLD
```

(1 row)

All arguments are specified in order. The result is upper case since \texttt{uppercase} is specified as \texttt{true}. Another example is:
Chapter 4. SQL Syntax

SELECT concat_lower_or_upper('Hello', 'World');

- -----------------------
  hello world
(1 row)

Here, the uppercase parameter is omitted, so it receives its default value of false, resulting in lower case output. In positional notation, arguments can be omitted from right to left so long as they have defaults.

4.3.2. Using Named Notation

In named notation, each argument’s name is specified using = to separate it from the argument expression. For example:

SELECT concat_lower_or_upper(a = 'Hello', b = 'World');

- -----------------------
  hello world
(1 row)

Again, the argument uppercase was omitted so it is set to false implicitly. One advantage of using named notation is that the arguments may be specified in any order, for example:

SELECT concat_lower_or_upper(a = 'Hello', b = 'World', uppercase = true);

- -----------------------
  HELLO WORLD
(1 row)

SELECT concat_lower_or_upper(a = 'Hello', uppercase = true, b = 'World');

- -----------------------
  HELLO WORLD
(1 row)

An older syntax based on ":=" is supported for backward compatibility:

SELECT concat_lower_or_upper(a := 'Hello', uppercase := true, b := 'World');

- -----------------------
  HELLO WORLD
(1 row)

4.3.3. Using Mixed Notation

The mixed notation combines positional and named notation. However, as already mentioned, named arguments cannot precede positional arguments. For example:

SELECT concat_lower_or_upper('Hello', 'World', uppercase => true);

- -----------------------
In the above query, the arguments \texttt{a} and \texttt{b} are specified positionally, while \texttt{uppercase} is specified by name. In this example, that adds little except documentation. With a more complex function having numerous parameters that have default values, named or mixed notation can save a great deal of writing and reduce chances for error.

\textbf{Note:} Named and mixed call notations currently cannot be used when calling an aggregate function (but they do work when an aggregate function is used as a window function).
Chapter 5. Data Definition

This chapter covers how one creates the database structures that will hold one’s data. In a relational database, the raw data is stored in tables, so the majority of this chapter is devoted to explaining how tables are created and modified and what features are available to control what data is stored in the tables. Subsequently, we discuss how tables can be organized into schemas, and how privileges can be assigned to tables. Finally, we will briefly look at other features that affect the data storage, such as inheritance, views, functions, and triggers.

5.1. Table Basics

A table in a relational database is much like a table on paper: It consists of rows and columns. The number and order of the columns is fixed, and each column has a name. The number of rows is variable — it reflects how much data is stored at a given moment. SQL does not make any guarantees about the order of the rows in a table. When a table is read, the rows will appear in an unspecified order, unless sorting is explicitly requested. This is covered in Chapter 7. Furthermore, SQL does not assign unique identifiers to rows, so it is possible to have several completely identical rows in a table. This is a consequence of the mathematical model that underlies SQL but is usually not desirable. Later in this chapter we will see how to deal with this issue.

Each column has a data type. The data type constrains the set of possible values that can be assigned to a column and assigns semantics to the data stored in the column so that it can be used for computations. For instance, a column declared to be of a numerical type will not accept arbitrary text strings, and the data stored in such a column can be used for mathematical computations. By contrast, a column declared to be of a character string type will accept almost any kind of data but it does not lend itself to mathematical calculations, although other operations such as string concatenation are available.

PostgreSQL includes a sizable set of built-in data types that fit many applications. Users can also define their own data types. Most built-in data types have obvious names and semantics, so we defer a detailed explanation to Chapter 8. Some of the frequently used data types are integer for whole numbers, numeric for possibly fractional numbers, text for character strings, date for dates, time for time-of-day values, and timestamp for values containing both date and time.

To create a table, you use the aptly named CREATE TABLE command. In this command you specify at least a name for the new table, the names of the columns and the data type of each column. For example:

```sql
CREATE TABLE my_first_table (  
    first_column text,  
    second_column integer  
);  
```

This creates a table named `my_first_table` with two columns. The first column is named `first_column` and has a data type of `text`; the second column has the name `second_column` and the type `integer`. The table and column names follow the identifier syntax explained in Section 4.1.1. The type names are usually also identifiers, but there are some exceptions. Note that the column list is comma-separated and surrounded by parentheses.

Of course, the previous example was heavily contrived. Normally, you would give names to your tables and columns that convey what kind of data they store. So let’s look at a more realistic example:

```sql
CREATE TABLE products {  
```
Chapter 5. Data Definition

```
product_no integer,
name text,
price numeric
);
```

(The numeric type can store fractional components, as would be typical of monetary amounts.)

**Tip:** When you create many interrelated tables it is wise to choose a consistent naming pattern for the tables and columns. For instance, there is a choice of using singular or plural nouns for table names, both of which are favored by some theorist or other.

There is a limit on how many columns a table can contain. Depending on the column types, it is between 250 and 1600. However, defining a table with anywhere near this many columns is highly unusual and often a questionable design.

If you no longer need a table, you can remove it using the DROP TABLE command. For example:

```
DROP TABLE my_first_table;
DROP TABLE products;
```

Attempting to drop a table that does not exist is an error. Nevertheless, it is common in SQL script files to unconditionally try to drop each table before creating it, ignoring any error messages, so that the script works whether or not the table exists. (If you like, you can use the `DROP TABLE IF EXISTS` variant to avoid the error messages, but this is not standard SQL.)

If you need to modify a table that already exists, see Section 5.5 later in this chapter.

With the tools discussed so far you can create fully functional tables. The remainder of this chapter is concerned with adding features to the table definition to ensure data integrity, security, or convenience. If you are eager to fill your tables with data now you can skip ahead to Chapter 6 and read the rest of this chapter later.

### 5.2. Default Values

A column can be assigned a default value. When a new row is created and no values are specified for some of the columns, those columns will be filled with their respective default values. A data manipulation command can also request explicitly that a column be set to its default value, without having to know what that value is. (Details about data manipulation commands are in Chapter 6.)

If no default value is declared explicitly, the default value is the null value. This usually makes sense because a null value can be considered to represent unknown data.

In a table definition, default values are listed after the column data type. For example:

```
CREATE TABLE products {
    product_no integer,
    name text,
    price numeric DEFAULT 9.99
};
```

The default value can be an expression, which will be evaluated whenever the default value is inserted (*not* when the table is created). A common example is for a `timestamp` column to have a default of

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51
CURRENT_TIMESTAMP, so that it gets set to the time of row insertion. Another common example is generating a “serial number” for each row. In PostgreSQL this is typically done by something like:

```sql
CREATE TABLE products {
    product_no integer DEFAULT nextval('products_product_no_seq'),
    ...
};
```

where the `nextval()` function supplies successive values from a sequence object (see Section 9.16). This arrangement is sufficiently common that there’s a special shorthand for it:

```sql
CREATE TABLE products {
    product_no SERIAL,
    ...
};
```

The `SERIAL` shorthand is discussed further in Section 8.1.4.

### 5.3. Constraints

Data types are a way to limit the kind of data that can be stored in a table. For many applications, however, the constraint they provide is too coarse. For example, a column containing a product price should probably only accept positive values. But there is no standard data type that accepts only positive numbers. Another issue is that you might want to constrain column data with respect to other columns or rows. For example, in a table containing product information, there should be only one row for each product number.

To that end, SQL allows you to define constraints on columns and tables. Constraints give you as much control over the data in your tables as you wish. If a user attempts to store data in a column that would violate a constraint, an error is raised. This applies even if the value came from the default value definition.

#### 5.3.1. Check Constraints

A check constraint is the most generic constraint type. It allows you to specify that the value in a certain column must satisfy a Boolean (truth-value) expression. For instance, to require positive product prices, you could use:

```sql
CREATE TABLE products {
    product_no integer,
    name text,
    price numeric CHECK (price > 0)
};
```

As you see, the constraint definition comes after the data type, just like default value definitions. Default values and constraints can be listed in any order. A check constraint consists of the key word `CHECK` followed by an expression in parentheses. The check constraint expression should involve the column thus constrained, otherwise the constraint would not make too much sense.

You can also give the constraint a separate name. This clarifies error messages and allows you to refer to the constraint when you need to change it. The syntax is:
Chapter 5. Data Definition

CREATE TABLE products (  
  product_no integer,  
  name text,  
  price numeric CONSTRAINT positive_price CHECK (price > 0)  
);  

So, to specify a named constraint, use the key word CONSTRAINT followed by an identifier followed by the constraint definition. (If you don’t specify a constraint name in this way, the system chooses a name for you.)

A check constraint can also refer to several columns. Say you store a regular price and a discounted price, and you want to ensure that the discounted price is lower than the regular price:

CREATE TABLE products (  
  product_no integer,  
  name text,  
  price numeric CHECK (price > 0),  
  discounted_price numeric CHECK (discounted_price > 0),  
  CHECK (price > discounted_price)  
);  

The first two constraints should look familiar. The third one uses a new syntax. It is not attached to a particular column, instead it appears as a separate item in the comma-separated column list. Column definitions and these constraint definitions can be listed in mixed order.

We say that the first two constraints are column constraints, whereas the third one is a table constraint because it is written separately from any one column definition. Column constraints can also be written as table constraints, while the reverse is not necessarily possible, since a column constraint is supposed to refer to only the column it is attached to. (PostgreSQL doesn’t enforce that rule, but you should follow it if you want your table definitions to work with other database systems.) The above example could also be written as:

CREATE TABLE products (  
  product_no integer,  
  name text,  
  price numeric,  
  CHECK (price > 0),  
  discounted_price numeric,  
  CHECK (discounted_price > 0),  
  CHECK (price > discounted_price)  
);  

or even:

CREATE TABLE products (  
  product_no integer,  
  name text,  
  price numeric CHECK (price > 0),  
  discounted_price numeric,  
  CHECK (discounted_price > 0 AND price > discounted_price)  
);  

It’s a matter of taste.

Names can be assigned to table constraints in the same way as column constraints:

CREATE TABLE products {  
  product_no integer,  
  name text,  
  price numeric,  
  ...
Chapter 5. Data Definition

CREATE TABLE products {
    product_no integer,
    name text,
    price numeric,
    CHECK (price > 0),
    discounted_price numeric,
    CHECK (discounted_price > 0),
    CONSTRAINT valid_discount CHECK (price > discounted_price)
};

It should be noted that a check constraint is satisfied if the check expression evaluates to true or the null value. Since most expressions will evaluate to the null value if any operand is null, they will not prevent null values in the constrained columns. To ensure that a column does not contain null values, the not-null constraint described in the next section can be used.

Note: PostgreSQL does not support CHECK constraints that reference table data other than the new or updated row being checked. While a CHECK constraint that violates this rule may appear to work in simple tests, it cannot guarantee that the database will not reach a state in which the constraint condition is false (due to subsequent changes of the other row(s) involved). This would cause a database dump and reload to fail. The reload could fail even when the complete database state is consistent with the constraint, due to rows not being loaded in an order that will satisfy the constraint. If possible, use UNIQUE, EXCLUDE, or FOREIGN KEY constraints to express cross-row and cross-table restrictions.

If what you desire is a one-time check against other rows at row insertion, rather than a continuously-maintained consistency guarantee, a custom trigger can be used to implement that. (This approach avoids the dump/reload problem because pg_dump does not reinstall triggers until after reloading data, so that the check will not be enforced during a dump/reload.)

Note: PostgreSQL assumes that CHECK constraints’ conditions are immutable, that is, they will always give the same result for the same input row. This assumption is what justifies examining CHECK constraints only when rows are inserted or updated, and not at other times. (The warning above about not referencing other table data is really a special case of this restriction.)

An example of a common way to break this assumption is to reference a user-defined function in a CHECK expression, and then change the behavior of that function. PostgreSQL does not disallow that, but it will not notice if there are rows in the table that now violate the CHECK constraint. That would cause a subsequent database dump and reload to fail. The recommended way to handle such a change is to drop the constraint (using ALTER TABLE), adjust the function definition, and re-add the constraint, thereby rechecking it against all table rows.

5.3.2. Not-Null Constraints

A not-null constraint simply specifies that a column must not assume the null value. A syntax example:

CREATE TABLE products {
    product_no integer NOT NULL,
    name text NOT NULL,
    price numeric
};
A not-null constraint is always written as a column constraint. A not-null constraint is functionally equivalent to creating a check constraint `CHECK (column_name IS NOT NULL)`, but in PostgreSQL creating an explicit not-null constraint is more efficient. The drawback is that you cannot give explicit names to not-null constraints created this way.

Of course, a column can have more than one constraint. Just write the constraints one after another:

```sql
CREATE TABLE products (
    product_no integer NOT NULL,
    name text NOT NULL,
    price numeric NOT NULL CHECK (price > 0)
);
```

The order doesn’t matter. It does not necessarily determine in which order the constraints are checked.

The `NOT NULL` constraint has an inverse: the `NULL` constraint. This does not mean that the column must be null, which would surely be useless. Instead, this simply selects the default behavior that the column might be null. The `NULL` constraint is not present in the SQL standard and should not be used in portable applications. (It was only added to PostgreSQL to be compatible with some other database systems.) Some users, however, like it because it makes it easy to toggle the constraint in a script file. For example, you could start with:

```sql
CREATE TABLE products {
    product_no integer NULL,
    name text NULL,
    price numeric NULL
};
```

and then insert the `NOT` keyword where desired.

**Tip:** In most database designs the majority of columns should be marked not null.

### 5.3.3. Unique Constraints

Unique constraints ensure that the data contained in a column, or a group of columns, is unique among all the rows in the table. The syntax is:

```sql
CREATE TABLE products {
    product_no integer UNIQUE,
    name text,
    price numeric
};
```

when written as a column constraint, and:

```sql
CREATE TABLE products {
    product_no integer,
    name text,
    price numeric,
    UNIQUE (product_no)
};
```

when written as a table constraint.
To define a unique constraint for a group of columns, write it as a table constraint with the column
names separated by commas:

```
CREATE TABLE example (
    a integer,
    b integer,
    c integer,
    UNIQUE (a, c)
);
```

This specifies that the combination of values in the indicated columns is unique across the whole
table, though any one of the columns need not be (and ordinarily isn’t) unique.

You can assign your own name for a unique constraint, in the usual way:

```
CREATE TABLE products (
    product_no integer CONSTRAINT must_be_different UNIQUE,
    name text,
    price numeric
);
```

Adding a unique constraint will automatically create a unique B-tree index on the column or group of
columns listed in the constraint. A uniqueness restriction covering only some rows cannot be written
as a unique constraint, but it is possible to enforce such a restriction by creating a unique partial index.

In general, a unique constraint is violated if there is more than one row in the table where the values of
all of the columns included in the constraint are equal. However, two null values are never considered
equal in this comparison. That means even in the presence of a unique constraint it is possible to
store duplicate rows that contain a null value in at least one of the constrained columns. This behavior
conforms to the SQL standard, but we have heard that other SQL databases might not follow this rule.
So be careful when developing applications that are intended to be portable.

### 5.3.4. Primary Keys

A primary key constraint indicates that a column, or group of columns, can be used as a unique
identifier for rows in the table. This requires that the values be both unique and not null. So, the
following two table definitions accept the same data:

```
CREATE TABLE products (
    product_no integer UNIQUE NOT NULL,
    name text,
    price numeric
);
```

```
CREATE TABLE products (
    product_no integer PRIMARY KEY,
    name text,
    price numeric
);
```

Primary keys can span more than one column; the syntax is similar to unique constraints:

```
CREATE TABLE example (  
    ```
Chapter 5. Data Definition

Adding a primary key will automatically create a unique B-tree index on the column or group of columns listed in the primary key, and will force the column(s) to be marked \textit{NOT NULL}.

A table can have at most one primary key. (There can be any number of unique and not-null constraints, which are functionally almost the same thing, but only one can be identified as the primary key.) Relational database theory dictates that every table must have a primary key. This rule is not enforced by PostgreSQL, but it is usually best to follow it.

Primary keys are useful both for documentation purposes and for client applications. For example, a GUI application that allows modifying row values probably needs to know the primary key of a table to be able to identify rows uniquely. There are also various ways in which the database system makes use of a primary key if one has been declared; for example, the primary key defines the default target column(s) for foreign keys referencing its table.

5.3.5. Foreign Keys

A foreign key constraint specifies that the values in a column (or a group of columns) must match the values appearing in some row of another table. We say this maintains the \textit{referential integrity} between two related tables.

Say you have the product table that we have used several times already:

\begin{verbatim}
CREATE TABLE products {
    product_no integer PRIMARY KEY,
    name text,
    price numeric
};
\end{verbatim}

Let’s also assume you have a table storing orders of those products. We want to ensure that the orders table only contains orders of products that actually exist. So we define a foreign key constraint in the orders table that references the products table:

\begin{verbatim}
CREATE TABLE orders {
    order_id integer PRIMARY KEY,
    product_no integer REFERENCES products (product_no),
    quantity integer
};
\end{verbatim}

Now it is impossible to create orders with non-NULL \texttt{product_no} entries that do not appear in the products table.

We say that in this situation the orders table is the \textit{referencing} table and the products table is the \textit{referenced} table. Similarly, there are referencing and referenced columns.

You can also shorten the above command to:

\begin{verbatim}
CREATE TABLE orders {
    order_id integer PRIMARY KEY,
    product_no integer REFERENCES products,
};
\end{verbatim}
because in absence of a column list the primary key of the referenced table is used as the referenced column(s).

A foreign key can also constrain and reference a group of columns. As usual, it then needs to be written in table constraint form. Here is a contrived syntax example:

```sql
CREATE TABLE t1 (
    a integer PRIMARY KEY,
    b integer,
    c integer,
    FOREIGN KEY (b, c) REFERENCES other_table (c1, c2)
);
```

Of course, the number and type of the constrained columns need to match the number and type of the referenced columns.

You can assign your own name for a foreign key constraint, in the usual way.

A table can have more than one foreign key constraint. This is used to implement many-to-many relationships between tables. Say you have tables about products and orders, but now you want to allow one order to contain possibly many products (which the structure above did not allow). You could use this table structure:

```sql
CREATE TABLE products {
    product_no integer PRIMARY KEY,
    name text,
    price numeric
};

CREATE TABLE orders {
    order_id integer PRIMARY KEY,
    shipping_address text,
    ...
};

CREATE TABLE order_items {
    product_no integer REFERENCES products,
    order_id integer REFERENCES orders,
    quantity integer,
    PRIMARY KEY (product_no, order_id)
};
```

Notice that the primary key overlaps with the foreign keys in the last table.
We know that the foreign keys disallow creation of orders that do not relate to any products. But what if a product is removed after an order is created that references it? SQL allows you to handle that as well. Intuitively, we have a few options:

- Disallow deleting a referenced product
- Delete the orders as well
- Something else?
To illustrate this, let’s implement the following policy on the many-to-many relationship example above: when someone wants to remove a product that is still referenced by an order (via \texttt{order\_items}), we disallow it. If someone removes an order, the order items are removed as well:

\begin{verbatim}
CREATE TABLE products {
    product_no integer PRIMARY KEY,
    name text,
    price numeric
};

CREATE TABLE orders {
    order_id integer PRIMARY KEY,
    shipping_address text,
    ...
};

CREATE TABLE order\_items (
    product_no integer REFERENCES products ON DELETE RESTRICT,
    order_id integer REFERENCES orders ON DELETE CASCADE,
    quantity integer,
    PRIMARY KEY (product_no, order_id)
);
\end{verbatim}

Restricting and cascading deletes are the two most common options. \texttt{RESTRICT} prevents deletion of a referenced row. \texttt{NO ACTION} means that if any referencing rows still exist when the constraint is checked, an error is raised; this is the default behavior if you do not specify anything. (The essential difference between these two choices is that \texttt{NO ACTION} allows the check to be deferred until later in the transaction, whereas \texttt{RESTRICT} does not.) \texttt{CASCADE} specifies that when a referenced row is deleted, row(s) referencing it should be automatically deleted as well. There are two other options: \texttt{SET NULL} and \texttt{SET DEFAULT}. These cause the referencing column(s) in the referencing row(s) to be set to nulls or their default values, respectively, when the referenced row is deleted. Note that these do not excuse you from observing any constraints. For example, if an action specifies \texttt{SET DEFAULT} but the default value would not satisfy the foreign key constraint, the operation will fail.

Analogous to \texttt{ON DELETE} there is also \texttt{ON UPDATE} which is invoked when a referenced column is changed (updated). The possible actions are the same. In this case, \texttt{CASCADE} means that the updated values of the referenced column(s) should be copied into the referencing row(s).

Normally, a referencing row need not satisfy the foreign key constraint if any of its referencing columns are null. If \texttt{MATCH FULL} is added to the foreign key declaration, a referencing row escapes satisfying the constraint only if all its referencing columns are null (so a mix of null and non-null values is guaranteed to fail a \texttt{MATCH FULL} constraint). If you don’t want referencing rows to be able to avoid satisfying the foreign key constraint, declare the referencing column(s) as \texttt{NOT NULL}.

A foreign key must reference columns that either are a primary key or form a unique constraint. This means that the referenced columns always have an index (the one underlying the primary key or unique constraint); so checks on whether a referencing row has a match will be efficient. Since a \texttt{DELETE} of a row from the referenced table or an \texttt{UPDATE} of a referenced column will require a scan of the referencing table for rows matching the old value, it is often a good idea to index the referencing columns too. Because this is not always needed, and there are many choices available on how to index, declaration of a foreign key constraint does not automatically create an index on the referencing columns.

More information about updating and deleting data is in Chapter 6. Also see the description of foreign key constraint syntax in the reference documentation for \texttt{CREATE TABLE}. 

Chapter 5. Data Definition
5.3.6. Exclusion Constraints

Exclusion constraints ensure that if any two rows are compared on the specified columns or expressions using the specified operators, at least one of these operator comparisons will return false or null.

The syntax is:

```
CREATE TABLE circles (  
    c circle,  
    EXCLUDE USING gist (c WITH &&)  
);  
```

See also `CREATE TABLE ... CONSTRAINT ... EXCLUDE` for details.

Adding an exclusion constraint will automatically create an index of the type specified in the constraint declaration.

5.4. System Columns

Every table has several system columns that are implicitly defined by the system. Therefore, these names cannot be used as names of user-defined columns. (Note that these restrictions are separate from whether the name is a key word or not; quoting a name will not allow you to escape these restrictions.) You do not really need to be concerned about these columns; just know they exist.

**oid**

The object identifier (object ID) of a row. This column is only present if the table was created using `WITH OIDS`, or if the `default_with_oids` configuration variable was set at the time. This column is of type `oid` (same name as the column); see Section 8.18 for more information about the type.

**tableoid**

The OID of the table containing this row. This column is particularly handy for queries that select from inheritance hierarchies (see Section 5.9), since without it, it’s difficult to tell which individual table a row came from. The `tableoid` can be joined against the `oid` column of `pg_class` to obtain the table name.

**xmin**

The identity (transaction ID) of the inserting transaction for this row version. (A row version is an individual state of a row; each update of a row creates a new row version for the same logical row.)

**cmin**

The command identifier (starting at zero) within the inserting transaction.

**xmax**

The identity (transaction ID) of the deleting transaction, or zero for an undeleted row version. It is possible for this column to be nonzero in a visible row version. That usually indicates that the deleting transaction hasn’t committed yet, or that an attempted deletion was rolled back.

**cmax**

The command identifier within the deleting transaction, or zero.
The physical location of the row version within its table. Note that although the ctid can be used to locate the row version very quickly, a row's ctid will change if it is updated or moved by VACUUM FULL. Therefore ctid is useless as a long-term row identifier. The OID, or even better a user-defined serial number, should be used to identify logical rows.

OIDs are 32-bit quantities and are assigned from a single cluster-wide counter. In a large or long-lived database, it is possible for the counter to wrap around. Hence, it is bad practice to assume that OIDs are unique, unless you take steps to ensure that this is the case. If you need to identify the rows in a table, using a sequence generator is strongly recommended. However, OIDs can be used as well, provided that a few additional precautions are taken:

- A unique constraint should be created on the OID column of each table for which the OID will be used to identify rows. When such a unique constraint (or unique index) exists, the system takes care not to generate an OID matching an already-existing row. (Of course, this is only possible if the table contains fewer than $2^{32}$ (4 billion) rows, and in practice the table size had better be much less than that, or performance might suffer.)
- OIDs should never be assumed to be unique across tables; use the combination of tableoid and row OID if you need a database-wide identifier.
- Of course, the tables in question must be created WITH OIDS. As of PostgreSQL 8.1, WITHOUT OIDS is the default.

Transaction identifiers are also 32-bit quantities. In a long-lived database it is possible for transaction IDs to wrap around. This is not a fatal problem given appropriate maintenance procedures; see Chapter 24 for details. It is unwise, however, to depend on the uniqueness of transaction IDs over the long term (more than one billion transactions).

Command identifiers are also 32-bit quantities. This creates a hard limit of $2^{32}$ (4 billion) SQL commands within a single transaction. In practice this limit is not a problem — note that the limit is on the number of SQL commands, not the number of rows processed. Also, only commands that actually modify the database contents will consume a command identifier.

### 5.5. Modifying Tables

When you create a table and you realize that you made a mistake, or the requirements of the application change, you can drop the table and create it again. But this is not a convenient option if the table is already filled with data, or if the table is referenced by other database objects (for instance a foreign key constraint). Therefore PostgreSQL provides a family of commands to make modifications to existing tables. Note that this is conceptually distinct from altering the data contained in the table: here we are interested in altering the definition, or structure, of the table.

You can:

- Add columns
- Remove columns
- Add constraints
- Remove constraints
- Change default values
- Change column data types
- Rename columns
Chapter 5. Data Definition

- Rename tables

All these actions are performed using the ALTER TABLE command, whose reference page contains details beyond those given here.

### 5.5.1. Adding a Column

To add a column, use a command like:

```sql
ALTER TABLE products ADD COLUMN description text;
```

The new column is initially filled with whatever default value is given (null if you don’t specify a `DEFAULT` clause).

You can also define constraints on the column at the same time, using the usual syntax:

```sql
ALTER TABLE products ADD COLUMN description text CHECK (description <> '');
```

In fact all the options that can be applied to a column description in `CREATE TABLE` can be used here. Keep in mind however that the default value must satisfy the given constraints, or the `ADD` will fail. Alternatively, you can add constraints later (see below) after you’ve filled in the new column correctly.

**Tip:** Adding a column with a default requires updating each row of the table (to store the new column value). However, if no default is specified, PostgreSQL is able to avoid the physical update. So if you intend to fill the column with mostly nondefault values, it’s best to add the column with no default, insert the correct values using `UPDATE`, and then add any desired default as described below.

### 5.5.2. Removing a Column

To remove a column, use a command like:

```sql
ALTER TABLE products DROP COLUMN description;
```

Whatever data was in the column disappears. Table constraints involving the column are dropped, too. However, if the column is referenced by a foreign key constraint of another table, PostgreSQL will not silently drop that constraint. You can authorize dropping everything that depends on the column by adding `CASCADE`:

```sql
ALTER TABLE products DROP COLUMN description CASCADE;
```

See Section 5.13 for a description of the general mechanism behind this.

### 5.5.3. Adding a Constraint

To add a constraint, the table constraint syntax is used. For example:

```sql
ALTER TABLE products ADD CHECK (name <> ' ');  
ALTER TABLE products ADD CONSTRAINT some_name UNIQUE (product_no);  
ALTER TABLE products ADD FOREIGN KEY (product_group_id) REFERENCES product_groups;
```

To add a not-null constraint, which cannot be written as a table constraint, use this syntax:
ALTER TABLE products ALTER COLUMN product_no SET NOT NULL;

The constraint will be checked immediately, so the table data must satisfy the constraint before it can be added.

### 5.5.4. Removing a Constraint

To remove a constraint you need to know its name. If you gave it a name then that’s easy. Otherwise the system assigned a generated name, which you need to find out. The psql command `\d tablename` can be helpful here; other interfaces might also provide a way to inspect table details. Then the command is:

```
ALTER TABLE products DROP CONSTRAINT some_name;
```

(If you are dealing with a generated constraint name like `$2`, don’t forget that you’ll need to double-quote it to make it a valid identifier.)

As with dropping a column, you need to add `CASCADE` if you want to drop a constraint that something else depends on. An example is that a foreign key constraint depends on a unique or primary key constraint on the referenced column(s).

This works the same for all constraint types except not-null constraints. To drop a not null constraint use:

```
ALTER TABLE products ALTER COLUMN product_no DROP NOT NULL;
```

(Recall that not-null constraints do not have names.)

### 5.5.5. Changing a Column’s Default Value

To set a new default for a column, use a command like:

```
ALTER TABLE products ALTER COLUMN price SET DEFAULT 7.77;
```

Note that this doesn’t affect any existing rows in the table, it just changes the default for future `INSERT` commands.

To remove any default value, use:

```
ALTER TABLE products ALTER COLUMN price DROP DEFAULT;
```

This is effectively the same as setting the default to null. As a consequence, it is not an error to drop a default where one hadn’t been defined, because the default is implicitly the null value.

### 5.5.6. Changing a Column’s Data Type

To convert a column to a different data type, use a command like:

```
ALTER TABLE products ALTER COLUMN price TYPE numeric(10,2);
```
Chapter 5. Data Definition

This will succeed only if each existing entry in the column can be converted to the new type by an implicit cast. If a more complex conversion is needed, you can add a USING clause that specifies how to compute the new values from the old.

PostgreSQL will attempt to convert the column’s default value (if any) to the new type, as well as any constraints that involve the column. But these conversions might fail, or might produce surprising results. It’s often best to drop any constraints on the column before altering its type, and then add back suitably modified constraints afterwards.

5.5.7. Renaming a Column

To rename a column:

```sql
ALTER TABLE products RENAME COLUMN product_no TO product_number;
```

5.5.8. Renaming a Table

To rename a table:

```sql
ALTER TABLE products RENAME TO items;
```

5.6. Privileges

When an object is created, it is assigned an owner. The owner is normally the role that executed the creation statement. For most kinds of objects, the initial state is that only the owner (or a superuser) can do anything with the object. To allow other roles to use it, privileges must be granted.

There are different kinds of privileges: SELECT, INSERT, UPDATE, DELETE, TRUNCATE, REFERENCES, TRIGGER, CREATE, CONNECT, TEMPORARY, EXECUTE, and USAGE. The privileges applicable to a particular object vary depending on the object’s type (table, function, etc). For complete information on the different types of privileges supported by PostgreSQL, refer to the GRANT reference page. The following sections and chapters will also show you how those privileges are used.

The right to modify or destroy an object is always the privilege of the owner only.

An object can be assigned to a new owner with an ALTER command of the appropriate kind for the object, e.g. ALTER TABLE. Superusers can always do this; ordinary roles can only do it if they are both the current owner of the object (or a member of the owning role) and a member of the new owning role.

To assign privileges, the GRANT command is used. For example, if joe is an existing role, and accounts is an existing table, the privilege to update the table can be granted with:

```sql
GRANT UPDATE ON accounts TO joe;
```

Writing ALL in place of a specific privilege grants all privileges that are relevant for the object type.
Chapter 5. Data Definition

The special “role” name PUBLIC can be used to grant a privilege to every role on the system. Also, “group” roles can be set up to help manage privileges when there are many users of a database — for details see Chapter 21.

To revoke a privilege, use the fittingly named REVOKE command:

```
REVOKE ALL ON accounts FROM PUBLIC;
```

The special privileges of the object owner (i.e., the right to do DROP, GRANT, REVOKE, etc.) are always implicit in being the owner, and cannot be granted or revoked. But the object owner can choose to revoke their own ordinary privileges, for example to make a table read-only for themselves as well as others.

Ordinarily, only the object’s owner (or a superuser) can grant or revoke privileges on an object. However, it is possible to grant a privilege “with grant option”, which gives the recipient the right to grant it in turn to others. If the grant option is subsequently revoked then all who received the privilege from that recipient (directly or through a chain of grants) will lose the privilege. For details see the GRANT and REVOKE reference pages.

5.7. Row Security Policies

In addition to the SQL-standard privilege system available through GRANT, tables can have row security policies that restrict, on a per-user basis, which rows can be returned by normal queries or inserted, updated, or deleted by data modification commands. This feature is also known as Row-Level Security. By default, tables do not have any policies, so that if a user has access privileges to a table according to the SQL privilege system, all rows within it are equally available for querying or updating.

When row security is enabled on a table (with ALTER TABLE ... ENABLE ROW LEVEL SECURITY), all normal access to the table for selecting rows or modifying rows must be allowed by a row security policy. (However, the table’s owner is typically not subject to row security policies.) If no policy exists for the table, a default-deny policy is used, meaning that no rows are visible or can be modified. Operations that apply to the whole table, such as TRUNCATE and REFERENCES, are not subject to row security.

Row security policies can be specific to commands, or to roles, or to both. A policy can be specified to apply to ALL commands, or to SELECT, INSERT, UPDATE, or DELETE. Multiple roles can be assigned to a given policy, and normal role membership and inheritance rules apply.

To specify which rows are visible or modifiable according to a policy, an expression is required that returns a Boolean result. This expression will be evaluated for each row prior to any conditions or functions coming from the user’s query. (The only exceptions to this rule are leakproof functions, which are guaranteed to not leak information; the optimizer may choose to apply such functions ahead of the row-security check.) Rows for which the expression does not return true will not be processed. Separate expressions may be specified to provide independent control over the rows which are visible and the rows which are allowed to be modified. Policy expressions are run as part of the query and with the privileges of the user running the query, although security-definer functions can be used to access data not available to the calling user.

Superusers and roles with the BYPASSRLS attribute always bypass the row security system when accessing a table. Table owners normally bypass row security as well, though a table owner can choose to be subject to row security with ALTER TABLE ... FORCE ROW LEVEL SECURITY.

Enabling and disabling row security, as well as adding policies to a table, is always the privilege of the table owner only.
Policies are created using the CREATE POLICY command, altered using the ALTER POLICY command, and dropped using the DROP POLICY command. To enable and disable row security for a given table, use the ALTER TABLE command.

Each policy has a name and multiple policies can be defined for a table. As policies are table-specific, each policy for a table must have a unique name. Different tables may have policies with the same name.

When multiple policies apply to a given query, they are combined using OR, so that a row is accessible if any policy allows it. This is similar to the rule that a given role has the privileges of all roles that they are a member of.

As a simple example, here is how to create a policy on the account relation to allow only members of the managers role to access rows, and only rows of their accounts:

```sql
CREATE TABLE accounts (manager text, company text, contact_email text);
ALTER TABLE accounts ENABLE ROW LEVEL SECURITY;
CREATE POLICY account_managers ON accounts TO managers USING (manager = current_user);
```

The policy above implicitly provides a WITH CHECK clause identical to its USING clause, so that the constraint applies both to rows selected by a command (so a manager cannot SELECT, UPDATE, or DELETE existing rows belonging to a different manager) and to rows modified by a command (so rows belonging to a different manager cannot be created via INSERT or UPDATE).

If no role is specified, or the special user name PUBLIC is used, then the policy applies to all users on the system. To allow all users to access only their own row in a users table, a simple policy can be used:

```sql
CREATE POLICY user_policy ON users USING (user_name = current_user);
```

This works similarly to the previous example.

To use a different policy for rows that are being added to the table compared to those rows that are visible, multiple policies can be combined. This pair of policies would allow all users to view all rows in the users table, but only modify their own:

```sql
CREATE POLICY user_sel_policy ON users FOR SELECT USING (true);
CREATE POLICY user_mod_policy ON users USING (user_name = current_user);
```

In a SELECT command, these two policies are combined using OR, with the net effect being that all rows can be selected. In other command types, only the second policy applies, so that the effects are the same as before.

Row security can also be disabled with the ALTER TABLE command. Disabling row security does not remove any policies that are defined on the table; they are simply ignored. Then all rows in the table are visible and modifiable, subject to the standard SQL privileges system.

Below is a larger example of how this feature can be used in production environments. The table passwd emulates a Unix password file:

```sql
-- Simple passwd-file based example
CREATE TABLE passwd {
```
Chapter 5. Data Definition

CREATE ROLE admin; -- Administrator
CREATE ROLE bob; -- Normal user
CREATE ROLE alice; -- Normal user

-- Populate the table
INSERT INTO passwd VALUES
  ('admin','xxx',0,0,'Admin','111-222-3333',null,'/root','/bin/dash');
INSERT INTO passwd VALUES
  ('bob','xxx',1,1,'Bob','123-456-7890',null,'/home/bob','/bin/zsh');
INSERT INTO passwd VALUES
  ('alice','xxx',2,1,'Alice','098-765-4321',null,'/home/alice','/bin/zsh');

-- Be sure to enable row level security on the table
ALTER TABLE passwd ENABLE ROW LEVEL SECURITY;

-- Create policies
-- Administrator can see all rows and add any rows
CREATE POLICY admin_all ON passwd TO admin USING (true) WITH CHECK (true);
-- Normal users can view all rows
CREATE POLICY all_view ON passwd FOR SELECT USING (true);
-- Normal users can update their own records, but
-- limit which shells a normal user is allowed to set
CREATE POLICY user_mod ON passwd FOR UPDATE
  USING (current_user = user_name)
  WITH CHECK (current_user = user_name AND shell IN ('/bin/bash','/bin/sh','/bin/dash','/bin/zsh','/bin/tcsh'));

-- Allow admin all normal rights
GRANT SELECT, INSERT, UPDATE, DELETE ON passwd TO admin;
-- Users only get select access on public columns
GRANT SELECT
  (user_name, uid, gid, real_name, home_phone, extra_info, home_dir, shell)
ON passwd TO public;
-- Allow users to update certain columns
GRANT UPDATE
  (pwhash, real_name, home_phone, extra_info, shell)
ON passwd TO public;

As with any security settings, it’s important to test and ensure that the system is behaving as expected. Using the example above, this demonstrates that the permission system is working properly.

-- admin can view all rows and fields
postgres=> set role admin;
SET
Chapter 5. Data Definition

postgres=> table passwd;
user_name | pwhash | uid | gid | real_name | home_phone | extra_info | home_dir | shell
-----------+--------+-----+-----+-----------+--------------+------------+-------------+-----------
admin | xxx | 0 | 0 | Admin | 111-222-3333 | | /root | /bin/dash
bob | xxx | 1 | 1 | Bob | 123-456-7890 | | /home/bob | /bin/zsh
alice | xxx | 2 | 1 | Alice | 098-765-4321 | | /home/alice | /bin/zsh
(3 rows)

-- Test what Alice is able to do
postgres=> set role alice;
SET
postgres=> table passwd;
ERROR: permission denied for relation passwd
postgres=> select user_name,real_name,home_phone,extra_info,home_dir,shell from passwd;
user_name | real_name | home_phone | extra_info | home_dir | shell
-----------+-----------+--------------+------------+-------------+-----------
admin | Admin | 111-222-3333 | | /root | /bin/dash
bob | Bob | 123-456-7890 | | /home/bob | /bin/zsh
alice | Alice | 098-765-4321 | | /home/alice | /bin/zsh
(3 rows)

postgres=> update passwd set user_name = 'joe';
ERROR: permission denied for relation passwd
-- Alice is allowed to change her own real_name, but no others
postgres=> update passwd set real_name = 'Alice Doe';
UPDATE 1
postgres=> update passwd set real_name = 'John Doe' where user_name = 'admin';
UPDATE 0
postgres=> update passwd set shell = '/bin/xx';
ERROR: new row violates WITH CHECK OPTION for "passwd"
postgres=> delete from passwd;
ERROR: permission denied for relation passwd
postgres=> insert into passwd (user_name) values ('xxx');
ERROR: permission denied for relation passwd
-- Alice can change her own password; RLS silently prevents updating other rows
postgres=> update passwd set pwhash = 'abc';
UPDATE 1

Referential integrity checks, such as unique or primary key constraints and foreign key references, always bypass row security to ensure that data integrity is maintained. Care must be taken when developing schemas and row level policies to avoid “covert channel” leaks of information through such referential integrity checks.

In some contexts it is important to be sure that row security is not being applied. For example, when taking a backup, it could be disastrous if row security silently caused some rows to be omitted from the backup. In such a situation, you can set the row_security configuration parameter to off. This does not in itself bypass row security; what it does is throw an error if any query’s results would get filtered by a policy. The reason for the error can then be investigated and fixed.

In the examples above, the policy expressions consider only the current values in the row to be accessed or updated. This is the simplest and best-performing case; when possible, it’s best to design row security applications to work this way. If it is necessary to consult other rows or other tables to make a policy decision, that can be accomplished using sub-SELECTs, or functions that contain SELECTs, in the policy expressions. Be aware however that such accesses can create race conditions that could allow information leakage if care is not taken. As an example, consider the following table design:
-- definition of privilege groups
CREATE TABLE groups (group_id int PRIMARY KEY,
  group_name text NOT NULL);

INSERT INTO groups VALUES
  (1, 'low'),
  (2, 'medium'),
  (5, 'high');

GRANT ALL ON groups TO alice; -- alice is the administrator
GRANT SELECT ON groups TO public;

-- definition of users’ privilege levels
CREATE TABLE users (user_name text PRIMARY KEY,
  group_id int NOT NULL REFERENCES groups);

INSERT INTO users VALUES
  ('alice', 5),
  ('bob', 2),
  ('mallory', 2);

GRANT ALL ON users TO alice;
GRANT SELECT ON users TO public;

-- table holding the information to be protected
CREATE TABLE information (info text,
  group_id int NOT NULL REFERENCES groups);

INSERT INTO information VALUES
  ('barely secret', 1),
  ('slightly secret', 2),
  ('very secret', 5);

ALTER TABLE information ENABLE ROW LEVEL SECURITY;

-- a row should be visible to/updatable by users whose security group_id is
-- greater than or equal to the row’s group_id
CREATE POLICY fp_s ON information FOR SELECT
  USING (group_id < (SELECT group_id FROM users WHERE user_name = current_user));
CREATE POLICY fp_u ON information FOR UPDATE
  USING (group_id < (SELECT group_id FROM users WHERE user_name = current_user));

-- we rely only on RLS to protect the information table
GRANT ALL ON information TO public;

Now suppose that alice wishes to change the “slightly secret” information, but decides that mallory
should not be trusted with the new content of that row, so she does:

BEGIN;
UPDATE users SET group_id = 1 WHERE user_name = 'mallory';
UPDATE information SET info = 'secret from mallory' WHERE group_id = 2;
COMMIT;

That looks safe; there is no window wherein mallory should be able to see the “secret from mallory”
string. However, there is a race condition here. If mallory is concurrently doing, say,

SELECT * FROM information WHERE group_id = 2 FOR UPDATE;
and her transaction is in READ COMMITTED mode, it is possible for her to see “secret from mallory”. That happens if her transaction reaches the information row just after alice’s does. It blocks waiting for alice’s transaction to commit, then fetches the updated row contents thanks to the FOR UPDATE clause. However, it does not fetch an updated row for the implicit SELECT from users, because that sub-SELECT did not have FOR UPDATE; instead the users row is read with the snapshot taken at the start of the query. Therefore, the policy expression tests the old value of mallory’s privilege level and allows her to see the updated row.

There are several ways around this problem. One simple answer is to use SELECT ... FOR SHARE in sub-SELECTs in row security policies. However, that requires granting UPDATE privilege on the referenced table (here users) to the affected users, which might be undesirable. (But another row security policy could be applied to prevent them from actually exercising that privilege; or the sub-SELECT could be embedded into a security definer function.) Also, heavy concurrent use of row share locks on the referenced table could pose a performance problem, especially if updates of it are frequent. Another solution, practical if updates of the referenced table are infrequent, is to take an exclusive lock on the referenced table when updating it, so that no concurrent transactions could be examining old row values. Or one could just wait for all concurrent transactions to end after committing an update of the referenced table and before making changes that rely on the new security situation.

For additional details see CREATE POLICY and ALTER TABLE.

5.8. Schemas

A PostgreSQL database cluster contains one or more named databases. Users and groups of users are shared across the entire cluster, but no other data is shared across databases. Any given client connection to the server can access only the data in a single database, the one specified in the connection request.

**Note:** Users of a cluster do not necessarily have the privilege to access every database in the cluster. Sharing of user names means that there cannot be different users named, say, joe in two databases in the same cluster; but the system can be configured to allow joe access to only some of the databases.

A database contains one or more named schemas, which in turn contain tables. Schemas also contain other kinds of named objects, including data types, functions, and operators. The same object name can be used in different schemas without conflict; for example, both schema1 and myschema can contain tables named mytable. Unlike databases, schemas are not rigidly separated: a user can access objects in any of the schemas in the database they are connected to, if they have privileges to do so.

There are several reasons why one might want to use schemas:

- To allow many users to use one database without interfering with each other.
- To organize database objects into logical groups to make them more manageable.
- Third-party applications can be put into separate schemas so they do not collide with the names of other objects.

Schemas are analogous to directories at the operating system level, except that schemas cannot be nested.
5.8.1. Creating a Schema

To create a schema, use the CREATE SCHEMA command. Give the schema a name of your choice. For example:

```
CREATE SCHEMA myschema;
``` 

To create or access objects in a schema, write a qualified name consisting of the schema name and table name separated by a dot:

```
schema.table
```

This works anywhere a table name is expected, including the table modification commands and the data access commands discussed in the following chapters. (For brevity we will speak of tables only, but the same ideas apply to other kinds of named objects, such as types and functions.)

Actually, the even more general syntax

```
database.schema.table
```

can be used too, but at present this is just for pro forma compliance with the SQL standard. If you write a database name, it must be the same as the database you are connected to.

So to create a table in the new schema, use:

```
CREATE TABLE myschema.mytable ( ... );
```

To drop a schema if it’s empty (all objects in it have been dropped), use:

```
DROP SCHEMA myschema;
```

To drop a schema including all contained objects, use:

```
DROP SCHEMA myschema CASCADE;
```

See Section 5.13 for a description of the general mechanism behind this.

Often you will want to create a schema owned by someone else (since this is one of the ways to restrict the activities of your users to well-defined namespaces). The syntax for that is:

```
CREATE SCHEMA schema_name AUTHORIZATION user_name;
```

You can even omit the schema name, in which case the schema name will be the same as the user name. See Section 5.8.6 for how this can be useful.

Schema names beginning with `pg_` are reserved for system purposes and cannot be created by users.
5.8.2. The Public Schema

In the previous sections we created tables without specifying any schema names. By default such tables (and other objects) are automatically put into a schema named “public”. Every new database contains such a schema. Thus, the following are equivalent:

```sql
CREATE TABLE products ( ... );
```

and:

```sql
CREATE TABLE public.products ( ... );
```

5.8.3. The Schema Search Path

Qualified names are tedious to write, and it’s often best not to wire a particular schema name into applications anyway. Therefore tables are often referred to by unqualified names, which consist of just the table name. The system determines which table is meant by following a search path, which is a list of schemas to look in. The first matching table in the search path is taken to be the one wanted. If there is no match in the search path, an error is reported, even if matching table names exist in other schemas in the database.

The ability to create like-named objects in different schemas complicates writing a query that references precisely the same objects every time. It also opens up the potential for users to change the behavior of other users’ queries, maliciously or accidentally. Due to the prevalence of unqualified names in queries and their use in PostgreSQL internals, adding a schema to `search_path` effectively trusts all users having CREATE privilege on that schema. When you run an ordinary query, a malicious user able to create objects in a schema of your search path can take control and execute arbitrary SQL functions as though you executed them.

The first schema named in the search path is called the current schema. Aside from being the first schema searched, it is also the schema in which new tables will be created if the `CREATE TABLE` command does not specify a schema name.

To show the current search path, use the following command:

```sql
SHOW search_path;
```

In the default setup this returns:

```
search_path
------------
"$user", public
```

The first element specifies that a schema with the same name as the current user is to be searched. If no such schema exists, the entry is ignored. The second element refers to the public schema that we have seen already.

The first schema in the search path that exists is the default location for creating new objects. That is the reason that by default objects are created in the public schema. When objects are referenced in any other context without schema qualification (table modification, data modification, or query commands) the search path is traversed until a matching object is found. Therefore, in the default configuration, any unqualified access again can only refer to the public schema.

To put our new schema in the path, we use:
Chapter 5. Data Definition

```sql
SET search_path TO myschema,public;
```

(We omit the $user here because we have no immediate need for it.) And then we can access the table without schema qualification:

```sql
DROP TABLE mytable;
```

Also, since `myschema` is the first element in the path, new objects would by default be created in it. We could also have written:

```sql
SET search_path TO myschema;
```

Then we no longer have access to the public schema without explicit qualification. There is nothing special about the public schema except that it exists by default. It can be dropped, too.

See also Section 9.25 for other ways to manipulate the schema search path.

The search path works in the same way for data type names, function names, and operator names as it does for table names. Data type and function names can be qualified in exactly the same way as table names. If you need to write a qualified operator name in an expression, there is a special provision: you must write

```sql
OPERATOR(schema.operator)
```

This is needed to avoid syntactic ambiguity. An example is:

```sql
SELECT 3 OPERATOR(pg_catalog.+) 4;
```

In practice one usually relies on the search path for operators, so as not to have to write anything so ugly as that.

### 5.8.4. Schemas and Privileges

By default, users cannot access any objects in schemas they do not own. To allow that, the owner of the schema must grant the `USAGE` privilege on the schema. To allow users to make use of the objects in the schema, additional privileges might need to be granted, as appropriate for the object.

A user can also be allowed to create objects in someone else’s schema. To allow that, the `CREATE` privilege on the schema needs to be granted. Note that by default, everyone has `CREATE` and `USAGE` privileges on the schema `public`. This allows all users that are able to connect to a given database to create objects in its `public` schema. Some usage patterns call for revoking that privilege:

```sql
REVOKE CREATE ON SCHEMA public FROM PUBLIC;
```

(The first “public” is the schema, the second “public” means “every user”. In the first sense it is an identifier, in the second sense it is a key word, hence the different capitalization; recall the guidelines from Section 4.1.1.)

### 5.8.5. The System Catalog Schema

In addition to `public` and user-created schemas, each database contains a `pg_catalog` schema, which contains the system tables and all the built-in data types, functions, and operators. `pg_catalog` is always effectively part of the search path. If it is not named explicitly in the path then it is implicitly searched before searching the path’s schemas. This ensures that built-in names will always be findable.
However, you can explicitly place `pg_catalog` at the end of your search path if you prefer to have user-defined names override built-in names.

Since system table names begin with `pg_`, it is best to avoid such names to ensure that you won’t suffer a conflict if some future version defines a system table named the same as your table. (With the default search path, an unqualified reference to your table name would then be resolved as the system table instead.) System tables will continue to follow the convention of having names beginning with `pg_`, so that they will not conflict with unqualified user-table names so long as users avoid the `pg_` prefix.

### 5.8.6. Usage Patterns

Schemas can be used to organize your data in many ways. A **secure schema usage pattern** prevents untrusted users from changing the behavior of other users’ queries. When a database does not use a secure schema usage pattern, users wishing to securely query that database would take protective action at the beginning of each session. Specifically, they would begin each session by setting `search_path` to the empty string or otherwise removing non-superuser-writable schemas from `search_path`. There are a few usage patterns easily supported by the default configuration:

- **Constrain ordinary users to user-private schemas.** To implement this, issue `REVOKE CREATE ON SCHEMA public FROM PUBLIC`, and create a schema for each user with the same name as that user. Recall that the default search path starts with `$user`, which resolves to the user name. Therefore, if each user has a separate schema, they access their own schemas by default. After adopting this pattern in a database where untrusted users had already logged in, consider auditing the public schema for objects named like objects in schema `pg_catalog`. This pattern is a secure schema usage pattern unless an untrusted user is the database owner or holds the `CREATEROLE` privilege, in which case no secure schema usage pattern exists.

- **Remove the public schema from the default search path,** by modifying `postgresql.conf` or by issuing `ALTER ROLE ALL SET search_path = "$user"`. Everyone retains the ability to create objects in the public schema, but only qualified names will choose those objects. While qualified table references are fine, calls to functions in the public schema will be unsafe or unreliable. If you create functions or extensions in the public schema, use the first pattern instead. Otherwise, like the first pattern, this is secure unless an untrusted user is the database owner or holds the `CREATEROLE` privilege.

- **Keep the default.** All users access the public schema implicitly. This simulates the situation where schemas are not available at all, giving a smooth transition from the non-schema-aware world. However, this is never a secure pattern. It is acceptable only when the database has a single user or a few mutually-trusting users.

For any pattern, to install shared applications (tables to be used by everyone, additional functions provided by third parties, etc.), put them into separate schemas. Remember to grant appropriate privileges to allow the other users to access them. Users can then refer to these additional objects by qualifying the names with a schema name, or they can put the additional schemas into their search path, as they choose.

### 5.8.7. Portability

In the SQL standard, the notion of objects in the same schema being owned by different users does
Chapter 5. Data Definition

not exist. Moreover, some implementations do not allow you to create schemas that have a different name than their owner. In fact, the concepts of schema and user are nearly equivalent in a database system that implements only the basic schema support specified in the standard. Therefore, many users consider qualified names to really consist of user_name.table_name. This is how PostgreSQL will effectively behave if you create a per-user schema for every user.

Also, there is no concept of a public schema in the SQL standard. For maximum conformance to the standard, you should not use the public schema.

Of course, some SQL database systems might not implement schemas at all, or provide namespace support by allowing (possibly limited) cross-database access. If you need to work with those systems, then maximum portability would be achieved by not using schemas at all.

5.9. Inheritance

PostgreSQL implements table inheritance, which can be a useful tool for database designers. (SQL:1999 and later define a type inheritance feature, which differs in many respects from the features described here.)

Let’s start with an example: suppose we are trying to build a data model for cities. Each state has many cities, but only one capital. We want to be able to quickly retrieve the capital city for any particular state. This can be done by creating two tables, one for state capitals and one for cities that are not capitals. However, what happens when we want to ask for data about a city, regardless of whether it is a capital or not? The inheritance feature can help to resolve this problem. We define the capitals table so that it inherits from cities:

```sql
CREATE TABLE cities {
    name text,
    population float,
    altitude int -- in feet
};

CREATE TABLE capitals {
    state char(2)
} INHERITS (cities);
```

In this case, the capitals table inherits all the columns of its parent table, cities. State capitals also have an extra column, state, that shows their state.

In PostgreSQL, a table can inherit from zero or more other tables, and a query can reference either all rows of a table or all rows of a table plus all of its descendant tables. The latter behavior is the default. For example, the following query finds the names of all cities, including state capitals, that are located at an altitude over 500 feet:

```sql
SELECT name, altitude
FROM cities
WHERE altitude > 500;
```

Given the sample data from the PostgreSQL tutorial (see Section 2.1), this returns:

```
<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
<tr>
<td>Madison</td>
<td>845</td>
</tr>
</tbody>
</table>
```
On the other hand, the following query finds all the cities that are not state capitals and are situated at an altitude over 500 feet:

```
SELECT name, altitude
FROM ONLY cities
WHERE altitude > 500;
```

<table>
<thead>
<tr>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>Mariposa</td>
<td>1953</td>
</tr>
</tbody>
</table>

Here the `ONLY` keyword indicates that the query should apply only to `cities`, and not any tables below `cities` in the inheritance hierarchy. Many of the commands that we have already discussed — `SELECT`, `UPDATE` and `DELETE` — support the `ONLY` keyword.

You can also write the table name with a trailing `*` to explicitly specify that descendant tables are included:

```
SELECT name, altitude
FROM cities*
WHERE altitude > 500;
```

Writing `*` is not necessary, since this behavior is the default (unless you have changed the setting of the `sql_inheritance` configuration option). However writing `*` might be useful to emphasize that additional tables will be searched.

In some cases you might wish to know which table a particular row originated from. There is a system column called `tableoid` in each table which can tell you the originating table:

```
SELECT c.tableoid, c.name, c.altitude
FROM cities c
WHERE c.altitude > 500;
```

which returns:

<table>
<thead>
<tr>
<th>tableoid</th>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>139793</td>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>139793</td>
<td>Mariposa</td>
<td>1953</td>
</tr>
<tr>
<td>139798</td>
<td>Madison</td>
<td>845</td>
</tr>
</tbody>
</table>

(If you try to reproduce this example, you will probably get different numeric OIDs.) By doing a join with `pg_class` you can see the actual table names:

```
SELECT p.relname, c.name, c.altitude
FROM cities c, pg_class p
WHERE c.altitude > 500 AND c.tableoid = p.oid;
```

which returns:

<table>
<thead>
<tr>
<th>relname</th>
<th>name</th>
<th>altitude</th>
</tr>
</thead>
<tbody>
<tr>
<td>cities</td>
<td>Las Vegas</td>
<td>2174</td>
</tr>
<tr>
<td>cities</td>
<td>Mariposa</td>
<td>1953</td>
</tr>
</tbody>
</table>
Another way to get the same effect is to use the regclass pseudo-type, which will print the table OID symbolically:

```
SELECT c.tableoid::regclass, c.name, c.altitude
FROM cities c
WHERE c.altitude > 500;
```

Inheritance does not automatically propagate data from `INSERT` or `COPY` commands to other tables in the inheritance hierarchy. In our example, the following `INSERT` statement will fail:

```
INSERT INTO cities (name, population, altitude, state)
VALUES ('Albany', NULL, NULL, 'NY');
```

We might hope that the data would somehow be routed to the `capitals` table, but this does not happen: `INSERT` always inserts into exactly the table specified. In some cases it is possible to redirect the insertion using a rule (see Chapter 39). However that does not help for the above case because the `cities` table does not contain the column `state`, and so the command will be rejected before the rule can be applied.

All check constraints and not-null constraints on a parent table are automatically inherited by its children, unless explicitly specified otherwise with `NO INHERIT` clauses. Other types of constraints (unique, primary key, and foreign key constraints) are not inherited.

A table can inherit from more than one parent table, in which case it has the union of the columns defined by the parent tables. Any columns declared in the child table’s definition are added to these. If the same column name appears in multiple parent tables, or in both a parent table and the child’s definition, then these columns are “merged” so that there is only one such column in the child table. To be merged, columns must have the same data types, else an error is raised. Inheritable check constraints and not-null constraints are merged in a similar fashion. Thus, for example, a merged column will be marked not-null if any one of the column definitions it came from is marked not-null.

Check constraints are merged if they have the same name, and the merge will fail if their conditions are different.

Table inheritance is typically established when the child table is created, using the `INHERITS` clause of the `CREATE TABLE` statement. Alternatively, a table which is already defined in a compatible way can have a new parent relationship added, using the `INHERIT` variant of `ALTER TABLE`. To do this the new child table must already include columns with the same names and types as the columns of the parent. It must also include check constraints with the same names and check expressions as those of the parent. Similarly an inheritance link can be removed from a child using the `NO INHERIT` variant of `ALTER TABLE`. Dynamically adding and removing inheritance links like this can be useful when the inheritance relationship is being used for table partitioning (see Section 5.10).

One convenient way to create a compatible table that will later be made a new child is to use the `LIKE` clause in `CREATE TABLE`. This creates a new table with the same columns as the source table. If there are any `CHECK` constraints defined on the source table, the `INCLUDING CONSTRAINTS` option to `LIKE` should be specified, as the new child must have constraints matching the parent to be considered compatible.

A parent table cannot be dropped while any of its children remain. Neither can columns or check constraints of child tables be dropped or altered if they are inherited from any parent tables. If you wish to remove a table and all of its descendants, one easy way is to drop the parent table with the `CASCADE` option (see Section 5.13).
ALTER TABLE will propagate any changes in column data definitions and check constraints down the inheritance hierarchy. Again, dropping columns that are depended on by other tables is only possible when using the CASCADE option. ALTER TABLE follows the same rules for duplicate column merging and rejection that apply during CREATE TABLE.

Inherited queries perform access permission checks on the parent table only. Thus, for example, granting UPDATE permission on the cities table implies permission to update rows in the capitals table as well, when they are accessed through cities. This preserves the appearance that the data is (also) in the parent table. But the capitals table could not be updated directly without an additional grant. Two exceptions to this rule are TRUNCATE and LOCK TABLE, where permissions on the child tables are always checked, whether they are processed directly or recursively via those commands performed on the parent table.

In a similar way, the parent table’s row security policies (see Section 5.7) are applied to rows coming from child tables during an inherited query. A child table’s policies, if any, are applied only when it is the table explicitly named in the query; and in that case, any policies attached to its parent(s) are ignored.

Foreign tables (see Section 5.11) can also be part of inheritance hierarchies, either as parent or child tables, just as regular tables can be. If a foreign table is part of an inheritance hierarchy then any operations not supported by the foreign table are not supported on the whole hierarchy either.

5.9.1. Caveats

Note that not all SQL commands are able to work on inheritance hierarchies. Commands that are used for data querying, data modification, or schema modification (e.g., SELECT, UPDATE, DELETE, most variants of ALTER TABLE, but not INSERT or ALTER TABLE ... RENAME) typically default to including child tables and support the ONLY notation to exclude them. Commands that do database maintenance and tuning (e.g., REINDEX, VACUUM) typically only work on individual, physical tables and do not support recursing over inheritance hierarchies. The respective behavior of each individual command is documented in its reference page (Reference I, SQL Commands).

A serious limitation of the inheritance feature is that indexes (including unique constraints) and foreign key constraints only apply to single tables, not to their inheritance children. This is true on both the referencing and referenced sides of a foreign key constraint. Thus, in the terms of the above example:

• If we declared cities.name to be UNIQUE or a PRIMARY KEY, this would not stop the capitals table from having rows with names duplicating rows in cities. And those duplicate rows would by default show up in queries from cities. In fact, by default capitals would have no unique constraint at all, and so could contain multiple rows with the same name. You could add a unique constraint to capitals, but this would not prevent duplication compared to cities.

• Similarly, if we were to specify that cities.name REFERENCES some other table, this constraint would not automatically propagate to capitals. In this case you could work around it by manually adding the same REFERENCES constraint to capitals.

• Specifying that another table’s column REFERENCES cities(name) would allow the other table to contain city names, but not capital names. There is no good workaround for this case.

These deficiencies will probably be fixed in some future release, but in the meantime considerable care is needed in deciding whether inheritance is useful for your application.
5.10. Partitioning

PostgreSQL supports basic table partitioning. This section describes why and how to implement partitioning as part of your database design.

5.10.1. Overview

Partitioning refers to splitting what is logically one large table into smaller physical pieces. Partitioning can provide several benefits:

- Query performance can be improved dramatically in certain situations, particularly when most of the heavily accessed rows of the table are in a single partition or a small number of partitions. The partitioning substitutes for leading columns of indexes, reducing index size and making it more likely that the heavily-used parts of the indexes fit in memory.
- When queries or updates access a large percentage of a single partition, performance can be improved by taking advantage of sequential scan of that partition instead of using an index and random access reads scattered across the whole table.
- Bulk loads and deletes can be accomplished by adding or removing partitions, if that requirement is planned into the partitioning design. ALTER TABLE NO INHERIT and DROP TABLE are both far faster than a bulk operation. These commands also entirely avoid the VACUUM overhead caused by a bulk DELETE.
- Seldom-used data can be migrated to cheaper and slower storage media.

The benefits will normally be worthwhile only when a table would otherwise be very large. The exact point at which a table will benefit from partitioning depends on the application, although a rule of thumb is that the size of the table should exceed the physical memory of the database server.

Currently, PostgreSQL supports partitioning via table inheritance. Each partition must be created as a child table of a single parent table. The parent table itself is normally empty; it exists just to represent the entire data set. You should be familiar with inheritance (see Section 5.9) before attempting to set up partitioning.

The following forms of partitioning can be implemented in PostgreSQL:

Range Partitioning

The table is partitioned into “ranges” defined by a key column or set of columns, with no overlap between the ranges of values assigned to different partitions. For example one might partition by date ranges, or by ranges of identifiers for particular business objects.

List Partitioning

The table is partitioned by explicitly listing which key values appear in each partition.

5.10.2. Implementing Partitioning

To set up a partitioned table, do the following:

1. Create the “master” table, from which all of the partitions will inherit.
Chapter 5. Data Definition

This table will contain no data. Do not define any check constraints on this table, unless you intend them to be applied equally to all partitions. There is no point in defining any indexes or unique constraints on it, either.

2. Create several “child” tables that each inherit from the master table. Normally, these tables will not add any columns to the set inherited from the master.

We will refer to the child tables as partitions, though they are in every way normal PostgreSQL tables (or, possibly, foreign tables).

3. Add table constraints to the partition tables to define the allowed key values in each partition.

Typical examples would be:

\[
\begin{align*}
&\text{CHECK ( } x = 1 \text{ )} \\
&\text{CHECK ( } \text{county IN ( 'Oxfordshire', 'Buckinghamshire', 'Warwickshire' )} \\
&\text{CHECK ( } \text{outletID } \geq 100 \text{ AND outletID } < 200 \text{ )}
\end{align*}
\]

Ensure that the constraints guarantee that there is no overlap between the key values permitted in different partitions. A common mistake is to set up range constraints like:

\[
\begin{align*}
&\text{CHECK ( } \text{outletID BETWEEN 100 AND 200} \text{ )} \\
&\text{CHECK ( } \text{outletID BETWEEN 200 AND 300} \text{ )}
\end{align*}
\]

This is wrong since it is not clear which partition the key value 200 belongs in.

Note that there is no difference in syntax between range and list partitioning; those terms are descriptive only.

4. For each partition, create an index on the key column(s), as well as any other indexes you might want. (The key index is not strictly necessary, but in most scenarios it is helpful. If you intend the key values to be unique then you should always create a unique or primary-key constraint for each partition.)

5. Optionally, define a trigger or rule to redirect data inserted into the master table to the appropriate partition.

6. Ensure that the constraint_exclusion configuration parameter is not disabled in \texttt{postgresql.conf}. If it is, queries will not be optimized as desired.

For example, suppose we are constructing a database for a large ice cream company. The company measures peak temperatures every day as well as ice cream sales in each region. Conceptually, we want a table like:

```sql
CREATE TABLE measurement (  
    city_id int not null,  
    logdate date not null,  
    peaktemp int,  
    unitsales int  
);
```

We know that most queries will access just the last week’s, month’s or quarter’s data, since the main use of this table will be to prepare online reports for management. To reduce the amount of old data that needs to be stored, we decide to only keep the most recent 3 years worth of data. At the beginning of each month we will remove the oldest month’s data.

In this situation we can use partitioning to help us meet all of our different requirements for the measurements table. Following the steps outlined above, partitioning can be set up as follows:

1. The master table is the \texttt{measurement} table, declared exactly as above.

2. Next we create one partition for each active month:

```sql
CREATE TABLE measurement_y2006m02 ( ) INHERITS (measurement);
```
CREATE TABLE measurement_y2006m03 ( ) INHERITS (measurement);
...
CREATE TABLE measurement_y2006m02 ( ) INHERITS (measurement);
CREATE TABLE measurement_y2006m03 ( ) INHERITS (measurement);
CREATE TABLE measurement_y2006m04 ( ) INHERITS (measurement);
Each of the partitions are complete tables in their own right, but they inherit their definitions from
the measurement table.

This solves one of our problems: deleting old data. Each month, all we will need to do is perform
a DROP TABLE on the oldest child table and create a new child table for the new month’s data.

3. We must provide non-overlapping table constraints. Rather than just creating the partition tables
as above, the table creation script should really be:

CREATE TABLE measurement_y2006m02 (CHECK ( logdate >= DATE '2006-02-01' AND logdate < DATE '2006-03-01') )
INHERITS (measurement);
CREATE TABLE measurement_y2006m03 (CHECK ( logdate >= DATE '2006-03-01' AND logdate < DATE '2006-04-01') )
INHERITS (measurement);
...
CREATE TABLE measurement_y2007m11 (CHECK ( logdate >= DATE '2007-11-01' AND logdate < DATE '2007-12-01') )
INHERITS (measurement);
CREATE TABLE measurement_y2007m12 (CHECK ( logdate >= DATE '2007-12-01' AND logdate < DATE '2008-01-01') )
INHERITS (measurement);
CREATE TABLE measurement_y2008m01 (CHECK ( logdate >= DATE '2008-01-01' AND logdate < DATE '2008-02-01') )
INHERITS (measurement);

4. We probably need indexes on the key columns too:

CREATE INDEX measurement_y2006m02_logdate ON measurement_y2006m02 (logdate);
CREATE INDEX measurement_y2006m03_logdate ON measurement_y2006m03 (logdate);
...
CREATE INDEX measurement_y2007m11_logdate ON measurement_y2007m11 (logdate);
CREATE INDEX measurement_y2007m12_logdate ON measurement_y2007m12 (logdate);
CREATE INDEX measurement_y2008m01_logdate ON measurement_y2008m01 (logdate);

We choose not to add further indexes at this time.

5. We want our application to be able to say INSERT INTO measurement ... and have the data
be redirected into the appropriate partition table. We can arrange that by attaching a suitable
trigger function to the master table. If data will be added only to the latest partition, we can use a
very simple trigger function:

CREATE OR REPLACE FUNCTION measurement_insert_trigger()
RETURNS TRIGGER AS $$
BEGIN
    INSERT INTO measurement_y2008m01 VALUES (NEW.*);
    RETURN NULL;
END;
$$
LANGUAGE plpgsql;

After creating the function, we create a trigger which calls the trigger function:

CREATE TRIGGER insert_measurement_trigger
BEFORE INSERT ON measurement
FOR EACH ROW EXECUTE PROCEDURE measurement_insert_trigger();

We must redefine the trigger function each month so that it always points to the current partition.
The trigger definition does not need to be updated, however.
We might want to insert data and have the server automatically locate the partition into which the row should be added. We could do this with a more complex trigger function, for example:

```sql
CREATE OR REPLACE FUNCTION measurement_insert_trigger()
RETURNS TRIGGER AS $$
BEGIN
    IF ( NEW.logdate >= DATE '2006-02-01' AND
        NEW.logdate < DATE '2006-03-01' ) THEN
        INSERT INTO measurement_y2006m02 VALUES (NEW.*);
    ELSIF ( NEW.logdate >= DATE '2006-03-01' AND
            NEW.logdate < DATE '2006-04-01' ) THEN
        INSERT INTO measurement_y2006m03 VALUES (NEW.*);
    ... 
    ELSIF ( NEW.logdate >= DATE '2008-01-01' AND
            NEW.logdate < DATE '2008-02-01' ) THEN
        INSERT INTO measurement_y2008m01 VALUES (NEW.*);
    ELSE
        RAISE EXCEPTION 'Date out of range. Fix the measurement_insert_trigger() function!';
    END IF;
    RETURN NULL;
END;
$$
LANGUAGE plpgsql;
```

The trigger definition is the same as before. Note that each `IF` test must exactly match the `CHECK` constraint for its partition.

While this function is more complex than the single-month case, it doesn’t need to be updated as often, since branches can be added in advance of being needed.

**Note:** In practice it might be best to check the newest partition first, if most inserts go into that partition. For simplicity we have shown the trigger’s tests in the same order as in other parts of this example.

As we can see, a complex partitioning scheme could require a substantial amount of DDL. In the above example we would be creating a new partition each month, so it might be wise to write a script that generates the required DDL automatically.

### 5.10.3. Managing Partitions

Normally the set of partitions established when initially defining the table are not intended to remain static. It is common to want to remove old partitions of data and periodically add new partitions for new data. One of the most important advantages of partitioning is precisely that it allows this otherwise painful task to be executed nearly instantaneously by manipulating the partition structure, rather than physically moving large amounts of data around.

The simplest option for removing old data is simply to drop the partition that is no longer necessary:

```sql
DROP TABLE measurement_y2006m02;
```

This can very quickly delete millions of records because it doesn’t have to individually delete every record.
Another option that is often preferable is to remove the partition from the partitioned table but retain access to it as a table in its own right:

```
ALTER TABLE measurement_y2006m02 NO INHERIT measurement;
```

This allows further operations to be performed on the data before it is dropped. For example, this is often a useful time to back up the data using `COPY`, `pg_dump`, or similar tools. It might also be a useful time to aggregate data into smaller formats, perform other data manipulations, or run reports.

Similarly we can add a new partition to handle new data. We can create an empty partition in the partitioned table just as the original partitions were created above:

```
CREATE TABLE measurement_y2008m02
  (CHECK ( logdate >= DATE '2008-02-01' AND logdate < DATE '2008-03-01' ) ) INHERITS (measurement);
```

As an alternative, it is sometimes more convenient to create the new table outside the partition structure, and make it a proper partition later. This allows the data to be loaded, checked, and transformed prior to it appearing in the partitioned table:

```
CREATE TABLE measurement_y2008m02
  (LIKE measurement INCLUDING DEFAULTS INCLUDING CONSTRAINTS);
ALTER TABLE measurement_y2008m02 ADD CONSTRAINT y2008m02
  CHECK ( logdate >= DATE '2008-02-01' AND logdate < DATE '2008-03-01' ) ;
\copy measurement_y2008m02 from 'measurement_y2008m02'
-- possibly some other data preparation work
ALTER TABLE measurement_y2008m02 INHERIT measurement;
```

### 5.10.4. Partitioning and Constraint Exclusion

*Constraint exclusion* is a query optimization technique that improves performance for partitioned tables defined in the fashion described above. As an example:

```
SET constraint_exclusion = on;
SELECT count(*) FROM measurement WHERE logdate >= DATE '2008-01-01';
```

Without constraint exclusion, the above query would scan each of the partitions of the `measurement` table. With constraint exclusion enabled, the planner will examine the constraints of each partition and try to prove that the partition need not be scanned because it could not contain any rows meeting the query’s `WHERE` clause. When the planner can prove this, it excludes the partition from the query plan.

You can use the `EXPLAIN` command to show the difference between a plan with `constraint_exclusion` on and a plan with it off. A typical unoptimized plan for this type of table setup is:

```
SET constraint_exclusion = off;
EXPLAIN SELECT count(*) FROM measurement WHERE logdate >= DATE '2008-01-01';
```

```
QUERY PLAN
-------------------------------------------------------------------------------
Aggregate (cost=158.66..158.68 rows=1 width=0)
  ->  Append (cost=0.00..151.88 rows=2715 width=0)
    ->  Seq Scan on measurement (cost=0.00..30.38 rows=543 width=0)
```

83
Chapter 5. Data Definition

Some or all of the partitions might use index scans instead of full-table sequential scans, but the point here is that there is no need to scan the older partitions at all to answer this query. When we enable constraint exclusion, we get a significantly cheaper plan that will deliver the same answer:

```
SET constraint_exclusion = on;
EXPLAIN SELECT count(*) FROM measurement WHERE logdate >='2008-01-01';
```

```
QUERY PLAN
-----------------------------------------------------------------------------------------------
Aggregate (cost=63.47..63.48 rows=1 width=0)  
  ->  Append (cost=0.00..60.75 rows=1086 width=0)  
    ->  Seq Scan on measurement (cost=0.00..30.38 rows=543 width=0)  
      Filter: (logdate >='2008-01-01')  
    ->  Seq Scan on measurement_y2008m01 measurement (cost=0.00..30.38 rows=543 width=0)  
      Filter: (logdate >='2008-01-01')
```

Note that constraint exclusion is driven only by CHECK constraints, not by the presence of indexes. Therefore it isn’t necessary to define indexes on the key columns. Whether an index needs to be created for a given partition depends on whether you expect that queries that scan the partition will generally scan a large part of the partition or just a small part. An index will be helpful in the latter case but not the former.

The default (and recommended) setting of constraint_exclusion is actually neither on nor off, but an intermediate setting called partition, which causes the technique to be applied only to queries that are likely to be working on partitioned tables. The on setting causes the planner to examine CHECK constraints in all queries, even simple ones that are unlikely to benefit.

5.10.5. Alternative Partitioning Methods

A different approach to redirecting inserts into the appropriate partition table is to set up rules, instead of a trigger, on the master table. For example:

```
CREATE RULE measurement_insert_y2006m02 AS
ON INSERT TO measurement WHERE
  ( logdate >= '2006-02-01' AND logdate < '2006-03-01' )
DO INSTEAD
  INSERT INTO measurement_y2006m02 VALUES (NEW.*);
```

```
CREATE RULE measurement_insert_y2008m01 AS
ON INSERT TO measurement WHERE
  ( logdate >= '2008-01-01' AND logdate < '2008-02-01' )
DO INSTEAD
  INSERT INTO measurement_y2008m01 VALUES (NEW.*);
```
A rule has significantly more overhead than a trigger, but the overhead is paid once per query rather than once per row, so this method might be advantageous for bulk-insert situations. In most cases, however, the trigger method will offer better performance.

Be aware that \texttt{COPY} ignores rules. If you want to use \texttt{COPY} to insert data, you’ll need to copy into the correct partition table rather than into the master. \texttt{COPY} does fire triggers, so you can use it normally if you use the trigger approach.

Another disadvantage of the rule approach is that there is no simple way to force an error if the set of rules doesn’t cover the insertion date; the data will silently go into the master table instead.

Partitioning can also be arranged using a \texttt{UNION ALL} view, instead of table inheritance. For example,

\begin{verbatim}
CREATE VIEW measurement AS
    SELECT * FROM measurement_y2006m02
UNION ALL SELECT * FROM measurement_y2006m03
... UNION ALL SELECT * FROM measurement_y2007m11
UNION ALL SELECT * FROM measurement_y2007m12
UNION ALL SELECT * FROM measurement_y2008m01;
\end{verbatim}

However, the need to recreate the view adds an extra step to adding and dropping individual partitions of the data set. In practice this method has little to recommend it compared to using inheritance.

\section*{5.10.6. Caveats}

The following caveats apply to partitioned tables:

\begin{itemize}
\item There is no automatic way to verify that all of the \texttt{CHECK} constraints are mutually exclusive. It is safer to create code that generates partitions and creates and/or modifies associated objects than to write each by hand.
\item The schemes shown here assume that the partition key column(s) of a row never change, or at least do not change enough to require it to move to another partition. An \texttt{UPDATE} that attempts to do that will fail because of the \texttt{CHECK} constraints. If you need to handle such cases, you can put suitable update triggers on the partition tables, but it makes management of the structure much more complicated.
\item If you are using manual \texttt{VACUUM} or \texttt{ANALYZE} commands, don’t forget that you need to run them on each partition individually. A command like:
\begin{verbatim}
ANALYZE measurement;
\end{verbatim}
will only process the master table.
\item \texttt{INSERT} statements with \texttt{ON CONFLICT} clauses are unlikely to work as expected, as the \texttt{ON CONFLICT} action is only taken in case of unique violations on the specified target relation, not its child relations.
\end{itemize}

The following caveats apply to constraint exclusion:

\begin{itemize}
\item Constraint exclusion only works when the query’s \texttt{WHERE} clause contains constants (or externally supplied parameters). For example, a comparison against a non-immutable function such as \texttt{CURRENT_TIMESTAMP} cannot be optimized, since the planner cannot know which partition the function value might fall into at run time.
\end{itemize}
• Keep the partitioning constraints simple, else the planner may not be able to prove that partitions don’t need to be visited. Use simple equality conditions for list partitioning, or simple range tests for range partitioning, as illustrated in the preceding examples. A good rule of thumb is that partitioning constraints should contain only comparisons of the partitioning column(s) to constants using B-tree-indexable operators.

• All constraints on all partitions of the master table are examined during constraint exclusion, so large numbers of partitions are likely to increase query planning time considerably. Partitioning using these techniques will work well with up to perhaps a hundred partitions; don’t try to use many thousands of partitions.

5.11. Foreign Data

PostgreSQL implements portions of the SQL/MED specification, allowing you to access data that resides outside PostgreSQL using regular SQL queries. Such data is referred to as foreign data. (Note that this usage is not to be confused with foreign keys, which are a type of constraint within the database.)

Foreign data is accessed with help from a foreign data wrapper. A foreign data wrapper is a library that can communicate with an external data source, hiding the details of connecting to the data source and obtaining data from it. There are some foreign data wrappers available as contrib modules; see Appendix F. Other kinds of foreign data wrappers might be found as third party products. If none of the existing foreign data wrappers suit your needs, you can write your own; see Chapter 55.

To access foreign data, you need to create a foreign server object, which defines how to connect to a particular external data source according to the set of options used by its supporting foreign data wrapper. Then you need to create one or more foreign tables, which define the structure of the remote data. A foreign table can be used in queries just like a normal table, but a foreign table has no storage in the PostgreSQL server. Whenever it is used, PostgreSQL asks the foreign data wrapper to fetch data from the external source, or transmit data to the external source in the case of update commands.

Accessing remote data may require authenticating to the external data source. This information can be provided by a user mapping, which can provide additional data such as user names and passwords based on the current PostgreSQL role.

For additional information, see CREATE FOREIGN DATA WRAPPER, CREATE SERVER, CREATE USER MAPPING, CREATE FOREIGN TABLE, and IMPORT FOREIGN SCHEMA.

5.12. Other Database Objects

Tables are the central objects in a relational database structure, because they hold your data. But they are not the only objects that exist in a database. Many other kinds of objects can be created to make the use and management of the data more efficient or convenient. They are not discussed in this chapter, but we give you a list here so that you are aware of what is possible:

• Views
• Functions and operators
• Data types and domains
Chapter 5. Data Definition

- Triggers and rewrite rules

Detailed information on these topics appears in Part V.

5.13. Dependency Tracking

When you create complex database structures involving many tables with foreign key constraints, views, triggers, functions, etc. you implicitly create a net of dependencies between the objects. For instance, a table with a foreign key constraint depends on the table it references.

To ensure the integrity of the entire database structure, PostgreSQL makes sure that you cannot drop objects that other objects still depend on. For example, attempting to drop the products table we considered in Section 5.3.5, with the orders table depending on it, would result in an error message like this:

```
DROP TABLE products;
ERROR:  cannot drop table products because other objects depend on it
DETAIL:  constraint orders_product_no_fkey on table orders depends on table products
HINT:    Use DROP ... CASCADE to drop the dependent objects too.
```

The error message contains a useful hint: if you do not want to bother deleting all the dependent objects individually, you can run:

```
DROP TABLE products CASCADE;
```

and all the dependent objects will be removed, as will any objects that depend on them, recursively.

In this case, it doesn’t remove the orders table, it only removes the foreign key constraint. It stops there because nothing depends on the foreign key constraint. (If you want to check what DROP ... CASCADE will do, run DROP without CASCADE and read the DETAIL output.)

Almost all DROP commands in PostgreSQL support specifying CASCADE. Of course, the nature of the possible dependencies varies with the type of the object. You can also write RESTRICT instead of CASCADE to get the default behavior, which is to prevent dropping objects that any other objects depend on.

Note: According to the SQL standard, specifying either RESTRICT or CASCADE is required in a DROP command. No database system actually enforces that rule, but whether the default behavior is RESTRICT or CASCADE varies across systems.

If a DROP command lists multiple objects, CASCADE is only required when there are dependencies outside the specified group. For example, when saying DROP TABLE tab1, tab2 the existence of a foreign key referencing tab1 from tab2 would not mean that CASCADE is needed to succeed.

For user-defined functions, PostgreSQL tracks dependencies associated with a function’s externally-visible properties, such as its argument and result types, but not dependencies that could only be known by examining the function body. As an example, consider this situation:

```
CREATE TYPE rainbow AS ENUM ('red', 'orange', 'yellow', 'green', 'blue', 'purple');
CREATE TABLE my_colors (color rainbow, note text);
CREATE FUNCTION get_color_note (rainbow) RETURNS text AS
```
‘SELECT note FROM my_colors WHERE color = $1’
LANGUAGE SQL;

(See Section 36.4 for an explanation of SQL-language functions.) PostgreSQL will be aware that the get_color_note function depends on the rainbow type: dropping the type would force dropping the function, because its argument type would no longer be defined. But PostgreSQL will not consider get_color_note to depend on the my_colors table, and so will not drop the function if the table is dropped. While there are disadvantages to this approach, there are also benefits. The function is still valid in some sense if the table is missing, though executing it would cause an error; creating a new table of the same name would allow the function to work again.
Chapter 6. Data Manipulation

The previous chapter discussed how to create tables and other structures to hold your data. Now it is time to fill the tables with data. This chapter covers how to insert, update, and delete table data. The chapter after this will finally explain how to extract your long-lost data from the database.

6.1. Inserting Data

When a table is created, it contains no data. The first thing to do before a database can be of much use is to insert data. Data is conceptually inserted one row at a time. Of course you can also insert more than one row, but there is no way to insert less than one row. Even if you know only some column values, a complete row must be created.

To create a new row, use the INSERT command. The command requires the table name and column values. For example, consider the products table from Chapter 5:

```sql
CREATE TABLE products (  
    product_no integer,  
    name text,  
    price numeric  
);
```

An example command to insert a row would be:

```sql
INSERT INTO products VALUES (1, 'Cheese', 9.99);
```

The data values are listed in the order in which the columns appear in the table, separated by commas. Usually, the data values will be literals (constants), but scalar expressions are also allowed.

The above syntax has the drawback that you need to know the order of the columns in the table. To avoid this you can also list the columns explicitly. For example, both of the following commands have the same effect as the one above:

```sql
INSERT INTO products (product_no, name, price) VALUES (1, 'Cheese', 9.99);
INSERT INTO products (name, price, product_no) VALUES ('Cheese', 9.99, 1);
```

Many users consider it good practice to always list the column names.

If you don’t have values for all the columns, you can omit some of them. In that case, the columns will be filled with their default values. For example:

```sql
INSERT INTO products (product_no, name) VALUES (1, 'Cheese');
INSERT INTO products VALUES (1, 'Cheese');
```

The second form is a PostgreSQL extension. It fills the columns from the left with as many values as are given, and the rest will be defaulted.

For clarity, you can also request default values explicitly, for individual columns or for the entire row:

```sql
INSERT INTO products (product_no, name, price) VALUES (1, 'Cheese', DEFAULT);
INSERT INTO products DEFAULT VALUES;
```

You can insert multiple rows in a single command:

```sql
INSERT INTO products (product_no, name, price) VALUES
```
(1, 'Cheese', 9.99),
(2, 'Bread', 1.99),
(3, 'Milk', 2.99);

It is also possible to insert the result of a query (which might be no rows, one row, or many rows):

```sql
INSERT INTO products (product_no, name, price)
SELECT product_no, name, price FROM new_products
WHERE release_date = 'today';
```

This provides the full power of the SQL query mechanism (Chapter 7) for computing the rows to be inserted.

**Tip:** When inserting a lot of data at the same time, consider using the COPY command. It is not as flexible as the INSERT command, but is more efficient. Refer to Section 14.4 for more information on improving bulk loading performance.

### 6.2. Updating Data

The modification of data that is already in the database is referred to as updating. You can update individual rows, all the rows in a table, or a subset of all rows. Each column can be updated separately; the other columns are not affected.

To update existing rows, use the UPDATE command. This requires three pieces of information:

1. The name of the table and column to update
2. The new value of the column
3. Which row(s) to update

Recall from Chapter 5 that SQL does not, in general, provide a unique identifier for rows. Therefore it is not always possible to directly specify which row to update. Instead, you specify which conditions a row must meet in order to be updated. Only if you have a primary key in the table (independent of whether you declared it or not) can you reliably address individual rows by choosing a condition that matches the primary key. Graphical database access tools rely on this fact to allow you to update rows individually.

For example, this command updates all products that have a price of 5 to have a price of 10:

```sql
UPDATE products SET price = 10 WHERE price = 5;
```

This might cause zero, one, or many rows to be updated. It is not an error to attempt an update that does not match any rows.

Let's look at that command in detail. First is the key word `UPDATE` followed by the table name. As usual, the table name can be schema-qualified, otherwise it is looked up in the path. Next is the key word `SET` followed by the column name, an equal sign, and the new column value. The new column value can be any scalar expression, not just a constant. For example, if you want to raise the price of all products by 10% you could use:

```sql
UPDATE products SET price = price * 1.10;
```
Chapter 6. Data Manipulation

As you see, the expression for the new value can refer to the existing value(s) in the row. We also left out the *WHERE* clause. If it is omitted, it means that all rows in the table are updated. If it is present, only those rows that match the *WHERE* condition are updated. Note that the equals sign in the *SET* clause is an assignment while the one in the *WHERE* clause is a comparison, but this does not create any ambiguity. Of course, the *WHERE* condition does not have to be an equality test. Many other operators are available (see Chapter 9). But the expression needs to evaluate to a Boolean result.

You can update more than one column in an *UPDATE* command by listing more than one assignment in the *SET* clause. For example:

```
UPDATE mytable SET a = 5, b = 3, c = 1 WHERE a > 0;
```

6.3. Deleting Data

So far we have explained how to add data to tables and how to change data. What remains is to discuss how to remove data that is no longer needed. Just as adding data is only possible in whole rows, you can only remove entire rows from a table. In the previous section we explained that SQL does not provide a way to directly address individual rows. Therefore, removing rows can only be done by specifying conditions that the rows to be removed have to match. If you have a primary key in the table then you can specify the exact row. But you can also remove groups of rows matching a condition, or you can remove all rows in the table at once.

You use the *DELETE* command to remove rows; the syntax is very similar to the *UPDATE* command. For instance, to remove all rows from the products table that have a price of 10, use:

```
DELETE FROM products WHERE price = 10;
```

If you simply write:

```
DELETE FROM products;
```

then all rows in the table will be deleted! Caveat programmer.

6.4. Returning Data From Modified Rows

Sometimes it is useful to obtain data from modified rows while they are being manipulated. The *INSERT, UPDATE, and DELETE* commands all have an optional *RETURNING* clause that supports this. Use of *RETURNING* avoids performing an extra database query to collect the data, and is especially valuable when it would otherwise be difficult to identify the modified rows reliably.

The allowed contents of a *RETURNING* clause are the same as a *SELECT* command’s output list (see Section 7.3). It can contain column names of the command’s target table, or value expressions using those columns. A common shorthand is *RETURNING *, which selects all columns of the target table in order.

In an *INSERT*, the data available to *RETURNING* is the row as it was inserted. This is not so useful in trivial inserts, since it would just repeat the data provided by the client. But it can be very handy when relying on computed default values. For example, when using a *serial* column to provide unique identifiers, *RETURNING* can return the ID assigned to a new row:
CREATE TABLE users (firstname text, lastname text, id serial primary key);

INSERT INTO users (firstname, lastname) VALUES ('Joe', 'Cool') RETURNING id;

The RETURNING clause is also very useful with INSERT ... SELECT.

In an UPDATE, the data available to RETURNING is the new content of the modified row. For example:

```
UPDATE products SET price = price * 1.10
    WHERE price <= 99.99
    RETURNING name, price AS new_price;
```

In a DELETE, the data available to RETURNING is the content of the deleted row. For example:

```
DELETE FROM products
    WHERE obsoletion_date = 'today'
    RETURNING *;
```

If there are triggers (Chapter 37) on the target table, the data available to RETURNING is the row as modified by the triggers. Thus, inspecting columns computed by triggers is another common use-case for RETURNING.
Chapter 7. Queries

The previous chapters explained how to create tables, how to fill them with data, and how to manipulate that data. Now we finally discuss how to retrieve the data from the database.

7.1. Overview

The process of retrieving or the command to retrieve data from a database is called a query. In SQL the SELECT command is used to specify queries. The general syntax of the SELECT command is

```
[WITH with_queries] SELECT select_list FROM table_expression [sort_specification]
```

The following sections describe the details of the select list, the table expression, and the sort specification. WITH queries are treated last since they are an advanced feature.

A simple kind of query has the form:

```
SELECT * FROM table1;
```

Assuming that there is a table called `table1`, this command would retrieve all rows and all user-defined columns from `table1`. (The method of retrieval depends on the client application. For example, the psql program will display an ASCII-art table on the screen, while client libraries will offer functions to extract individual values from the query result.) The select list specification `*` means all columns that the table expression happens to provide. A select list can also select a subset of the available columns or make calculations using the columns. For example, if `table1` has columns named `a`, `b`, and `c` (and perhaps others) you can make the following query:

```
SELECT a, b + c FROM table1;
```

(assuming that `b` and `c` are of a numerical data type). See Section 7.3 for more details.

```
FROM table1 is a simple kind of table expression: it reads just one table. In general, table expressions can be complex constructs of base tables, joins, and subqueries. But you can also omit the table expression entirely and use the SELECT command as a calculator:
```

```
SELECT 3 * 4;
```

This is more useful if the expressions in the select list return varying results. For example, you could call a function this way:

```
SELECT random();
```

7.2. Table Expressions

A `table expression` computes a table. The table expression contains a `FROM` clause that is optionally followed by `WHERE`, `GROUP BY`, and `HAVING` clauses. Trivial table expressions simply refer to a table on disk, a so-called base table, but more complex expressions can be used to modify or combine base tables in various ways.

The optional `WHERE`, `GROUP BY`, and `HAVING` clauses in the table expression specify a pipeline of successive transformations performed on the table derived in the `FROM` clause. All these transforma-
tions produce a virtual table that provides the rows that are passed to the select list to compute the output rows of the query.

### 7.2.1. The FROM Clause

The **FROM Clause** derives a table from one or more other tables given in a comma-separated table reference list.

```
FROM table_reference [, table_reference [, ...]]
```

A table reference can be a table name (possibly schema-qualified), or a derived table such as a sub-query, a JOIN construct, or complex combinations of these. If more than one table reference is listed in the FROM clause, the tables are cross-joined (that is, the Cartesian product of their rows is formed; see below). The result of the FROM list is an intermediate virtual table that can then be subject to transformations by the WHERE, GROUP BY, and HAVING clauses and is finally the result of the overall table expression.

When a table reference names a table that is the parent of a table inheritance hierarchy, the table reference produces rows of not only that table but all of its descendant tables, unless the key word **ONLY** precedes the table name. However, the reference produces only the columns that appear in the named table — any columns added in subtables are ignored.

Instead of writing **ONLY** before the table name, you can write *** after the table name to explicitly specify that descendant tables are included. Writing **_** is not necessary since that behavior is the default (unless you have changed the setting of the sql_inheritance configuration option). However writing **_** might be useful to emphasize that additional tables will be searched.

#### 7.2.1.1. Joined Tables

A joined table is a table derived from two other (real or derived) tables according to the rules of the particular join type. Inner, outer, and cross-joins are available. The general syntax of a joined table is

```
T1 join_type T2 [ join_condition ]
```

Joins of all types can be chained together, or nested: either or both T1 and T2 can be joined tables. Parentheses can be used around JOIN clauses to control the join order. In the absence of parentheses, JOIN clauses nest left-to-right.

**Join Types**

Cross join

```
T1 CROSS JOIN T2
```

For every possible combination of rows from T1 and T2 (i.e., a Cartesian product), the joined table will contain a row consisting of all columns in T1 followed by all columns in T2. If the tables have N and M rows respectively, the joined table will have N * M rows.

```
FROM T1 CROSS JOIN T2 is equivalent to FROM T1 INNER JOIN T2 ON TRUE (see below). It is also equivalent to FROM T1, T2.
```

**Note:** This latter equivalence does not hold exactly when more than two tables appear, because JOIN binds more tightly than comma. For example FROM T1 CROSS JOIN T2 INNER JOIN T3 ON condition is not the same as FROM T1, T2 INNER JOIN T3 ON condition because the condition can reference T1 in the first case but not the second.
Qualified joins

\[
T1 \{ [INNER] | \{ LEFT | RIGHT | FULL \} [OUTER] \} JOIN T2 \text{ ON boolean_expression}
\]

\[
T1 \{ [INNER] | \{ LEFT | RIGHT | FULL \} [OUTER] \} JOIN T2 \text{ USING ( join column list )}
\]

\[
T1 \text{ NATURAL \{ [INNER] | \{ LEFT | RIGHT | FULL \} [OUTER] \} JOIN T2}
\]

The words \texttt{INNER} and \texttt{OUTER} are optional in all forms. \texttt{INNER} is the default; \texttt{LEFT}, \texttt{RIGHT}, and \texttt{FULL} imply an outer join.

The \emph{join condition} is specified in the \texttt{ON} or \texttt{USING} clause, or implicitly by the word \texttt{NATURAL}. The join condition determines which rows from the two source tables are considered to “match”, as explained in detail below.

The possible types of qualified join are:

\textbf{INNER JOIN}

For each row \(R1\) of \(T1\), the joined table has a row for each row in \(T2\) that satisfies the join condition with \(R1\).

\textbf{LEFT OUTER JOIN}

First, an inner join is performed. Then, for each row in \(T1\) that does not satisfy the join condition with any row in \(T2\), a joined row is added with null values in columns of \(T2\). Thus, the joined table always has at least one row for each row in \(T1\).

\textbf{RIGHT OUTER JOIN}

First, an inner join is performed. Then, for each row in \(T2\) that does not satisfy the join condition with any row in \(T1\), a joined row is added with null values in columns of \(T1\). This is the converse of a left join: the result table will always have a row for each row in \(T2\).

\textbf{FULL OUTER JOIN}

First, an inner join is performed. Then, for each row in \(T1\) that does not satisfy the join condition with any row in \(T2\), a joined row is added with null values in columns of \(T2\). Also, for each row of \(T2\) that does not satisfy the join condition with any row in \(T1\), a joined row with null values in the columns of \(T1\) is added.

The \texttt{ON} clause is the most general kind of join condition: it takes a Boolean value expression of the same kind as is used in a \texttt{WHERE} clause. A pair of rows from \(T1\) and \(T2\) match if the \texttt{ON} expression evaluates to true.

The \texttt{USING} clause is a shorthand that allows you to take advantage of the specific situation where both sides of the join use the same name for the joining column(s). It takes a comma-separated list of the shared column names and forms a join condition that includes an equality comparison for each one. For example, joining \(T1\) and \(T2\) with \texttt{USING \{a, b\}} produces the join condition \(ON T1.a = T2.a \text{ AND } T1.b = T2.b\).

Furthermore, the output of \texttt{JOIN USING} suppresses redundant columns: there is no need to print both of the matched columns, since they must have equal values. While \texttt{JOIN ON} produces all columns from \(T1\) followed by all columns from \(T2\), \texttt{JOIN USING} produces one output column for each of the listed column pairs (in the listed order), followed by any remaining columns from \(T1\), followed by any remaining columns from \(T2\).

Finally, \texttt{NATURAL} is a shorthand form of \texttt{USING}: it forms a \texttt{USING} list consisting of all column names that appear in both input tables. As with \texttt{USING}, these columns appear only once in the
output table. If there are no common column names, \texttt{NATURAL \ JOIN} behaves like \texttt{JOIN ... ON TRUE}, producing a cross-product join.

\textbf{Note:} \texttt{USING} is reasonably safe from column changes in the joined relations since only the listed columns are combined. \texttt{NATURAL} is considerably more risky since any schema changes to either relation that cause a new matching column name to be present will cause the join to combine that new column as well.

To put this together, assume we have tables \texttt{t1}:

\begin{verbatim}
num | name
-----+------
 1   | a
 2   | b
 3   | c
\end{verbatim}

and \texttt{t2}:

\begin{verbatim}
num | value
-----+------
 1   | xxx
 3   | yyy
 5   | zzz
\end{verbatim}

then we get the following results for the various joins:

\begin{verbatim}
=> SELECT * FROM t1 CROSS JOIN t2;
 num | name | num | value
-----+------|-----+------
 1   | a    | 1   | xxx
 1   | a    | 3   | yyy
 1   | a    | 5   | zzz
 2   | b    | 1   | xxx
 2   | b    | 3   | yyy
 2   | b    | 5   | zzz
 3   | c    | 1   | xxx
 3   | c    | 3   | yyy
 3   | c    | 5   | zzz
(9 rows)

=> SELECT * FROM t1 INNER JOIN t2 ON t1.num = t2.num;
 num | name | num | value
-----+------|-----+------
 1   | a    | 1   | xxx
 3   | c    | 3   | yyy
(2 rows)

=> SELECT * FROM t1 INNER JOIN t2 USING (num);
 num | name | value
-----+------|------
 1   | a    | xxx
 3   | c    | yyy
(2 rows)
\end{verbatim}
Chapter 7. Queries

-> SELECT * FROM t1 NATURAL INNER JOIN t2;
  num  | name | value
  ------+------+-------
  1 | a | xxx
  3 | c | yyy
(2 rows)

-> SELECT * FROM t1 LEFT JOIN t2 ON t1.num = t2.num;
  num  | name | num | value
  ------+------+-----+-------
  1 | a | 1 | xxx
  2 | b | | |
  3 | c | 3 | yyy
(3 rows)

-> SELECT * FROM t1 LEFT JOIN t2 USING (num);
  num  | name | value
  ------+------+-------
  1 | a | xxx
  2 | b | |
  3 | c | yyy
(3 rows)

-> SELECT * FROM t1 RIGHT JOIN t2 ON t1.num = t2.num;
  num  | name | num | value
  ------+------+-----+-------
  1 | a | 1 | xxx
  3 | c | 3 | yyy
  | | 5 | zzz
(3 rows)

-> SELECT * FROM t1 FULL JOIN t2 ON t1.num = t2.num;
  num  | name | num | value
  ------+------+-----+-------
  1 | a | 1 | xxx
  2 | b | |
  3 | c | 3 | yyy
  | | 5 | zzz
(4 rows)

The join condition specified with ON can also contain conditions that do not relate directly to the join. This can prove useful for some queries but needs to be thought out carefully. For example:

-> SELECT * FROM t1 LEFT JOIN t2 ON t1.num = t2.num AND t2.value = 'xxx';
  num  | name | num | value
  ------+------+-----+-------
  1 | a | 1 | xxx
  2 | b ||
  3 | c | 3 | yyy
  | | 5 | zzz
(3 rows)

Notice that placing the restriction in the WHERE clause produces a different result:

-> SELECT * FROM t1 LEFT JOIN t2 ON t1.num = t2.num WHERE t2.value = 'xxx';
  num  | name | num | value
  ------+------+-----+-------
Chapter 7. Queries

This is because a restriction placed in the ON clause is processed before the join, while a restriction placed in the WHERE clause is processed after the join. That does not matter with inner joins, but it matters a lot with outer joins.

7.2.1.2. Table and Column Aliases

A temporary name can be given to tables and complex table references to be used for references to the derived table in the rest of the query. This is called a table alias.

To create a table alias, write

FROM table_reference AS alias

or

FROM table_reference alias

The AS key word is optional noise. alias can be any identifier.

A typical application of table aliases is to assign short identifiers to long table names to keep the join clauses readable. For example:

SELECT * FROM some_very_long_table_name s JOIN another_fairly_long_name a ON s.id = a.num;

The alias becomes the new name of the table reference so far as the current query is concerned — it is not allowed to refer to the table by the original name elsewhere in the query. Thus, this is not valid:

SELECT * FROM my_table AS m WHERE my_table.a > 5; -- wrong

Table aliases are mainly for notational convenience, but it is necessary to use them when joining a table to itself, e.g.:

SELECT * FROM people AS mother JOIN people AS child ON mother.id = child.mother_id;

Additionally, an alias is required if the table reference is a subquery (see Section 7.2.1.3).

Parentheses are used to resolve ambiguities. In the following example, the first statement assigns the alias b to the second instance of my_table, but the second statement assigns the alias to the result of the join:

SELECT * FROM my_table AS a CROSS JOIN my_table AS b ...
SELECT * FROM (my_table AS a CROSS JOIN my_table) AS b ...

Another form of table aliasing gives temporary names to the columns of the table, as well as the table itself:

FROM table_reference [AS] alias ( column1 [, column2 [, ...]] )

If fewer column aliases are specified than the actual table has columns, the remaining columns are not renamed. This syntax is especially useful for self-joins or subqueries.

98
When an alias is applied to the output of a JOIN clause, the alias hides the original name(s) within the JOIN. For example:

```
SELECT a.* FROM my_table AS a JOIN your_table AS b ON ...
```

is valid SQL, but:

```
SELECT a.* FROM (my_table AS a JOIN your_table AS b ON ...) AS c
```

is not valid; the table alias `a` is not visible outside the alias `c`.

### 7.2.1.3. Subqueries

Subqueries specifying a derived table must be enclosed in parentheses and must be assigned a table alias name (as in Section 7.2.1.2). For example:

```
FROM (SELECT * FROM table1) AS alias_name
```

This example is equivalent to `FROM table1 AS alias_name`. More interesting cases, which cannot be reduced to a plain join, arise when the subquery involves grouping or aggregation.

A subquery can also be a VALUES list:

```
FROM (VALUES ('anne', 'smith'), ('bob', 'jones'), ('joe', 'blow')) AS names(first, last)
```

Again, a table alias is required. Assigning alias names to the columns of the VALUES list is optional, but is good practice. For more information see Section 7.7.

### 7.2.1.4. Table Functions

Table functions are functions that produce a set of rows, made up of either base data types (scalar types) or composite data types (table rows). They are used like a table, view, or subquery in the FROM clause of a query. Columns returned by table functions can be included in SELECT, JOIN, or WHERE clauses in the same manner as columns of a table, view, or subquery.

Table functions may also be combined using the ROWS FROM syntax, with the results returned in parallel columns; the number of result rows in this case is that of the largest function result, with smaller results padded with null values to match.

```
function_call [WITH ORDINALITY] [[AS] table_alias [(column_alias [, ... ])]]
ROWS FROM( function_call [, ... ] ) [WITH ORDINALITY] [[AS] table_alias [(column_alias [, ... ])]]
```

If the WITH ORDINALITY clause is specified, an additional column of type bigint will be added to the function result columns. This column numbers the rows of the function result set, starting from 1. (This is a generalization of the SQL-standard syntax for UNNEST ... WITH ORDINALITY.) By default, the ordinal column is called `ordinality`, but a different column name can be assigned to it using an AS clause.

The special table function UNNEST may be called with any number of array parameters, and it returns a corresponding number of columns, as if UNNEST (Section 9.18) had been called on each parameter separately and combined using the ROWS FROM construct.

```
UNNEST( array_expression [, ... ] ) [WITH ORDINALITY] [[AS] table_alias [(column_alias [, ... ])]]
```
If no table_alias is specified, the function name is used as the table name; in the case of a ROWS FROM() construct, the first function’s name is used.

If column aliases are not supplied, then for a function returning a base data type, the column name is also the same as the function name. For a function returning a composite type, the result columns get the names of the individual attributes of the type.

Some examples:

CREATE TABLE foo (fooid int, foosubid int, fooname text);

CREATE FUNCTION getfoo(int) RETURNS SETOF foo AS $$
SELECT * FROM foo WHERE fooid = $1;
$$ LANGUAGE SQL;

SELECT * FROM getfoo(1) AS t1;

SELECT * FROM foo
WHERE foosubid IN (
    SELECT foosubid
    FROM getfoo(foo.fooid) z
    WHERE z.fooid = foo.fooid
);

CREATE VIEW vw_getfoo AS SELECT * FROM getfoo(1);

SELECT * FROM vw_getfoo;

In some cases it is useful to define table functions that can return different column sets depending on how they are invoked. To support this, the table function can be declared as returning the pseudotype record. When such a function is used in a query, the expected row structure must be specified in the query itself, so that the system can know how to parse and plan the query. This syntax looks like:

```
function_call [AS] alias (column_definition [, ... ])
function_call AS [alias] (column_definition [, ... ])
ROWS FROM( ... function_call AS (column_definition [, ... ])) [, ... ]
```

When not using the ROWS FROM() syntax, the column_definition list replaces the column alias list that could otherwise be attached to the FROM item; the names in the column definitions serve as column aliases. When using the ROWS FROM() syntax, a column_definition list can be attached to each member function separately; or if there is only one member function and no WITH ORDINALITY clause, a column_definition list can be written in place of a column alias list following ROWS FROM().

Consider this example:

```
SELECT *
    FROM dblink('dbname=mydb', 'SELECT proname, prosrc FROM pg_proc')
    AS t1(proname name, prosrc text)
    WHERE proname LIKE 'bytea%';
```

The dblink function (part of the dblink module) executes a remote query. It is declared to return record since it might be used for any kind of query. The actual column set must be specified in the calling query so that the parser knows, for example, what * should expand to.
Chapter 7. Queries

7.2.1.5. LATERAL Subqueries

Subqueries appearing in `FROM` can be preceded by the key word `LATERAL`. This allows them to reference columns provided by preceding `FROM` items. (Without `LATERAL`, each subquery is evaluated independently and so cannot cross-reference any other `FROM` item.)

Table functions appearing in `FROM` can also be preceded by the key word `LATERAL`, but for functions the key word is optional; the function’s arguments can contain references to columns provided by preceding `FROM` items in any case.

A `LATERAL` item can appear at top level in the `FROM` list, or within a `JOIN` tree. In the latter case it can also refer to any items that are on the left-hand side of a `JOIN` that it is on the right-hand side of.

When a `FROM` item contains `LATERAL` cross-references, evaluation proceeds as follows: for each row of the `FROM` item providing the cross-referenced column(s), or set of rows of multiple `FROM` items providing the columns, the `LATERAL` item is evaluated using that row or row set’s values of the columns. The resulting row(s) are joined as usual with the rows they were computed from. This is repeated for each row or set of rows from the column source table(s).

A trivial example of `LATERAL` is

```
SELECT * FROM foo, LATERAL (SELECT * FROM bar WHERE bar.id = foo.bar_id) ss;
```

This is not especially useful since it has exactly the same result as the more conventional

```
SELECT * FROM foo, bar WHERE bar.id = foo.bar_id;
```

`LATERAL` is primarily useful when the cross-referenced column is necessary for computing the row(s) to be joined. A common application is providing an argument value for a set-returning function. For example, supposing that `vertices(polygon)` returns the set of vertices of a polygon, we could identify close-together vertices of polygons stored in a table with:

```
SELECT p1.id, p2.id, v1, v2
FROM polygons p1, polygons p2,
     LATERAL vertices(p1.poly) v1,
     LATERAL vertices(p2.poly) v2
WHERE (v1 <-> v2) < 10 AND p1.id != p2.id;
```

This query could also be written

```
SELECT p1.id, p2.id, v1, v2
FROM polygons p1 CROSS JOIN LATERAL vertices(p1.poly) v1,
     polygons p2 CROSS JOIN LATERAL vertices(p2.poly) v2
WHERE (v1 <-> v2) < 10 AND p1.id != p2.id;
```

or in several other equivalent formulations. (As already mentioned, the `LATERAL` key word is unnecessary in this example, but we use it for clarity.)

It is often particularly handy to `LEFT JOIN` to a `LATERAL` subquery, so that source rows will appear in the result even if the `LATERAL` subquery produces no rows for them. For example, if `get_product_names()` returns the names of products made by a manufacturer, but some manufacturers in our table currently produce no products, we could find out which ones those are like this:

```
SELECT m.name
FROM manufacturers m LEFT JOIN LATERAL get_product_names(m.id) pname ON true
WHERE pname IS NULL;
```
7.2.2. The WHERE Clause

The syntax of the WHERE Clause is

```
WHERE search_condition
```

where `search_condition` is any value expression (see Section 4.2) that returns a value of type boolean.

After the processing of the FROM clause is done, each row of the derived virtual table is checked against the search condition. If the result of the condition is true, the row is kept in the output table, otherwise (i.e., if the result is false or null) it is discarded. The search condition typically references at least one column of the table generated in the FROM clause; this is not required, but otherwise the WHERE clause will be fairly useless.

**Note:** The join condition of an inner join can be written either in the WHERE clause or in the JOIN clause. For example, these table expressions are equivalent:

- `FROM a, b WHERE a.id = b.id AND b.val > 5`
- `FROM a INNER JOIN b ON (a.id = b.id) WHERE b.val > 5`
- or perhaps even:

```
FROM a NATURAL JOIN b WHERE b.val > 5
```

Which one of these you use is mainly a matter of style. The JOIN syntax in the FROM clause is probably not as portable to other SQL database management systems, even though it is in the SQL standard. For outer joins there is no choice: they must be done in the FROM clause. The ON or USING clause of an outer join is not equivalent to a WHERE condition, because it results in the addition of rows (for unmatched input rows) as well as the removal of rows in the final result.

Here are some examples of WHERE clauses:

```
SELECT ... FROM fdt WHERE c1 > 5
SELECT ... FROM fdt WHERE c1 IN (1, 2, 3)
SELECT ... FROM fdt WHERE c1 IN (SELECT c1 FROM t2)
SELECT ... FROM fdt WHERE c1 IN (SELECT c3 FROM t2 WHERE c2 = fdt.c1 + 10)
SELECT ... FROM fdt WHERE c1 BETWEEN (SELECT c3 FROM t2 WHERE c2 = fdt.c1 + 10) AND 100
SELECT ... FROM fdt WHERE EXISTS (SELECT c1 FROM t2 WHERE c2 > fdt.c1)
```

`fdt` is the table derived in the FROM clause. Rows that do not meet the search condition of the WHERE clause are eliminated from `fdt`. Notice the use of scalar subqueries as value expressions. Just like any other query, the subqueries can employ complex table expressions. Notice also how `fdt` is referenced in the subqueries. Qualifying `c1` as `fdt.c1` is only necessary if `c1` is also the name of a column in the derived input table of the subquery. But qualifying the column name adds clarity even when it is not needed. This example shows how the column naming scope of an outer query extends into its inner queries.
7.2.3. The `GROUP BY` and `HAVING` Clauses

After passing the `WHERE` filter, the derived input table might be subject to grouping, using the `GROUP BY` clause, and elimination of group rows using the `HAVING` clause.

```sql
SELECT select_list
FROM ...
[WHERE ...]
GROUP BY grouping_column_reference [, grouping_column_reference]...
```

The `GROUP BY` Clause is used to group together those rows in a table that have the same values in all the columns listed. The order in which the columns are listed does not matter. The effect is to combine each set of rows having common values into one group row that represents all rows in the group. This is done to eliminate redundancy in the output and/or compute aggregates that apply to these groups.

For instance:

```sql
=> SELECT * FROM test1;
 x | y
---+---
a | 3
c | 2
b | 5
a | 1
(4 rows)

=> SELECT x FROM test1 GROUP BY x;
 x
---
a
b
c
(3 rows)
```

In the second query, we could not have written `SELECT * FROM test1 GROUP BY x`, because there is no single value for the column `y` that could be associated with each group. The grouped-by columns can be referenced in the select list since they have a single value in each group.

In general, if a table is grouped, columns that are not listed in `GROUP BY` cannot be referenced except in aggregate expressions. An example with aggregate expressions is:

```sql
=> SELECT x, sum(y) FROM test1 GROUP BY x;
 x | sum
---+-----
a | 4
b | 5
c | 2
(3 rows)
```

Here `sum` is an aggregate function that computes a single value over the entire group. More information about the available aggregate functions can be found in Section 9.20.

**Tip:** Grouping without aggregate expressions effectively calculates the set of distinct values in a column. This can also be achieved using the `DISTINCT` clause (see Section 7.3.3).
Here is another example: it calculates the total sales for each product (rather than the total sales of all products):

```sql
SELECT product_id, p.name, (sum(s.units) * p.price) AS sales
FROM products p LEFT JOIN sales s USING (product_id)
GROUP BY product_id, p.name, p.price;
```

In this example, the columns `product_id`, `p.name`, and `p.price` must be in the `GROUP BY` clause since they are referenced in the query select list (but see below). The column `s.units` does not have to be in the `GROUP BY` list since it is only used in an aggregate expression (`sum(...)`), which represents the sales of a product. For each product, the query returns a summary row about all sales of the product.

If the products table is set up so that, say, `product_id` is the primary key, then it would be enough to group by `product_id` in the above example, since name and price would be functionally dependent on the product ID, and so there would be no ambiguity about which name and price value to return for each product ID group.

In strict SQL, `GROUP BY` can only group by columns of the source table but PostgreSQL extends this to also allow `GROUP BY` to group by columns in the select list. Grouping by value expressions instead of simple column names is also allowed.

If a table has been grouped using `GROUP BY`, but only certain groups are of interest, the `HAVING` clause can be used, much like a `WHERE` clause, to eliminate groups from the result. The syntax is:

```sql
SELECT select_list FROM ... [WHERE ...] GROUP BY ... HAVING boolean_expression
```

Expressions in the `HAVING` clause can refer both to grouped expressions and to ungrouped expressions (which necessarily involve an aggregate function).

Example:

```sql
=> SELECT x, sum(y) FROM test1 GROUP BY x HAVING sum(y) > 3;
   x | sum
  ---+-----
   a |  4
   b |  5
(2 rows)

=> SELECT x, sum(y) FROM test1 GROUP BY x HAVING x < 'c';
   x | sum
  ---+-----
   a |  4
   b |  5
(2 rows)
```

Again, a more realistic example:

```sql
SELECT product_id, p.name, (sum(s.units) * (p.price - p.cost)) AS profit
FROM products p LEFT JOIN sales s USING (product_id)
WHERE s.date > CURRENT_DATE - INTERVAL '4 weeks'
GROUP BY product_id, p.name, p.price, p.cost
HAVING sum(p.price * s.units) > 5000;
```

In the example above, the `WHERE` clause is selecting rows by a column that is not grouped (the expression is only true for sales during the last four weeks), while the `HAVING` clause restricts the output to
groups with total gross sales over 5000. Note that the aggregate expressions do not necessarily need to be the same in all parts of the query.

If a query contains aggregate function calls, but no `GROUP BY` clause, grouping still occurs: the result is a single group row (or perhaps no rows at all, if the single row is then eliminated by `HAVING`). The same is true if it contains a `HAVING` clause, even without any aggregate function calls or `GROUP BY` clause.

### 7.2.4. GROUPING SETS, CUBE, and ROLLUP

More complex grouping operations than those described above are possible using the concept of grouping sets. The data selected by the `FROM` and `WHERE` clauses is grouped separately by each specified grouping set, aggregates computed for each group just as for simple `GROUP BY` clauses, and then the results returned. For example:

```sql
-> SELECT * FROM items_sold;
brand | size | sales
-------+------+-------
Foo   | L    | 10
Foo   | M    | 20
Bar   | M    | 15
Bar   | L    | 5
(4 rows)

-> SELECT brand, size, sum(sales) FROM items_sold GROUP BY GROUPING SETS ((brand), (size), ());
brand | size | sum
-------+------+-----
Foo   |     | 30
Bar   |     | 20
| L    | 15
| M    | 35
|      | 50
(5 rows)
```

Each sublist of `GROUPING SETS` may specify zero or more columns or expressions and is interpreted the same way as though it were directly in the `GROUP BY` clause. An empty grouping set means that all rows are aggregated down to a single group (which is output even if no input rows were present), as described above for the case of aggregate functions with no `GROUP BY` clause.

References to the grouping columns or expressions are replaced by null values in result rows for grouping sets in which those columns do not appear. To distinguish which grouping a particular output row resulted from, see Table 9-55.

A shorthand notation is provided for specifying two common types of grouping set. A clause of the form

```
ROLLUP ( e1, e2, e3, ... )
```

represents the given list of expressions and all prefixes of the list including the empty list; thus it is equivalent to

```
GROUPING SETS (  
  ( e1, e2, e3, ... ),
  ...,  
  ( e1, e2 ),
```

105
This is commonly used for analysis over hierarchical data; e.g. total salary by department, division, and company-wide total.

A clause of the form

\[
\text{CUBE ( } e_1, e_2, \ldots \text{ )}
\]

represents the given list and all of its possible subsets (i.e. the power set). Thus

\[
\text{CUBE ( } a, b, c \text{ )}
\]

is equivalent to

\[
\text{GROUPING SETS (}
\begin{align*}
& ( a, b, c ), \\
& ( a, b ), \\
& ( a, c ), \\
& ( a ), \\
& ( b, c ), \\
& ( b ), \\
& ( c ), \\
& ( )
\end{align*}
\]

The individual elements of a CUBE or ROLLUP clause may be either individual expressions, or sublists of elements in parentheses. In the latter case, the sublists are treated as single units for the purposes of generating the individual grouping sets. For example:

\[
\text{CUBE ( } (a, b), \{c, d\} \text{ )}
\]

is equivalent to

\[
\text{GROUPING SETS (}
\begin{align*}
& ( a, b, c, d ), \\
& ( a, b ), \\
& ( c, d ), \\
& ( )
\end{align*}
\]

and

\[
\text{ROLLUP ( } a, (b, c), d \text{ )}
\]

is equivalent to

\[
\text{GROUPING SETS (}
\begin{align*}
& ( a, b, c, d ), \\
& ( a, b, c ), \\
& ( a ), \\
& ( )
\end{align*}
\]
Chapter 7. Queries

The CUBE and ROLLUP constructs can be used either directly in the GROUP BY clause, or nested inside a GROUPING SETS clause. If one GROUPING SETS clause is nested inside another, the effect is the same as if all the elements of the inner clause had been written directly in the outer clause.

If multiple grouping items are specified in a single GROUP BY clause, then the final list of grouping sets is the cross product of the individual items. For example:

GROUP BY a, CUBE (b, c), GROUPING SETS ((d), (e))

is equivalent to

GROUP BY GROUPING SETS (  
  (a, b, c, d), (a, b, c, e),  
  (a, b, d), (a, b, e),  
  (a, c, d), (a, c, e),  
  (a, d), (a, e)  
)

Note: The construct (a, b) is normally recognized in expressions as a row constructor. Within the GROUP BY clause, this does not apply at the top levels of expressions, and (a, b) is parsed as a list of expressions as described above. If for some reason you need a row constructor in a grouping expression, use ROW(a, b).

7.2.5. Window Function Processing

If the query contains any window functions (see Section 3.5, Section 9.21 and Section 4.2.8), these functions are evaluated after any grouping, aggregation, and HAVING filtering is performed. That is, if the query uses any aggregates, GROUP BY, or HAVING, then the rows seen by the window functions are the group rows instead of the original table rows from FROM/WHERE.

When multiple window functions are used, all the window functions having syntactically equivalent PARTITION BY and ORDER BY clauses in their window definitions are guaranteed to be evaluated in a single pass over the data. Therefore they will see the same sort ordering, even if the ORDER BY does not uniquely determine an ordering. However, no guarantees are made about the evaluation of functions having different PARTITION BY or ORDER BY specifications. (In such cases a sort step is typically required between the passes of window function evaluations, and the sort is not guaranteed to preserve ordering of rows that its ORDER BY sees as equivalent.)

Currently, window functions always require presorted data, and so the query output will be ordered according to one or another of the window functions’ PARTITION BY/ORDER BY clauses. It is not recommended to rely on this, however. Use an explicit top-level ORDER BY clause if you want to be sure the results are sorted in a particular way.

7.3. Select Lists

As shown in the previous section, the table expression in the SELECT command constructs an intermediate virtual table by possibly combining tables, views, eliminating rows, grouping, etc. This table
Chapter 7. Queries

is finally passed on to processing by the select list. The select list determines which columns of the intermediate table are actually output.

7.3.1. Select-List Items

The simplest kind of select list is * which emits all columns that the table expression produces. Otherwise, a select list is a comma-separated list of value expressions (as defined in Section 4.2). For instance, it could be a list of column names:

```
SELECT a, b, c FROM ...
```

The columns names a, b, and c are either the actual names of the columns of tables referenced in the FROM clause, or the aliases given to them as explained in Section 7.2.1.2. The name space available in the select list is the same as in the WHERE clause, unless grouping is used, in which case it is the same as in the HAVING clause.

If more than one table has a column of the same name, the table name must also be given, as in:

```
SELECT tbl1.a, tbl2.a, tbl1.b FROM ...
```

When working with multiple tables, it can also be useful to ask for all the columns of a particular table:

```
SELECT tbl1.*, tbl2.a FROM ...
```

See Section 8.16.5 for more about the table_name.* notation.

If an arbitrary value expression is used in the select list, it conceptually adds a new virtual column to the returned table. The value expression is evaluated once for each result row, with the row’s values substituted for any column references. But the expressions in the select list do not have to reference any columns in the table expression of the FROM clause; they can be constant arithmetic expressions, for instance.

7.3.2. Column Labels

The entries in the select list can be assigned names for subsequent processing, such as for use in an ORDER BY clause or for display by the client application. For example:

```
SELECT a AS value, b + c AS sum FROM ...
```

If no output column name is specified using AS, the system assigns a default column name. For simple column references, this is the name of the referenced column. For function calls, this is the name of the function. For complex expressions, the system will generate a generic name.

The AS keyword is optional, but only if the new column name does not match any PostgreSQL keyword (see Appendix C). To avoid an accidental match to a keyword, you can double-quote the column name. For example, VALUE is a keyword, so this does not work:

```
SELECT a value, b + c AS sum FROM ...
```

but this does:

```
SELECT a "value", b + c AS sum FROM ...
```
For protection against possible future keyword additions, it is recommended that you always either write AS or double-quote the output column name.

Note: The naming of output columns here is different from that done in the FROM clause (see Section 7.2.1.2). It is possible to rename the same column twice, but the name assigned in the select list is the one that will be passed on.

### 7.3.3. DISTINCT

After the select list has been processed, the result table can optionally be subject to the elimination of duplicate rows. The DISTINCT key word is written directly after SELECT to specify this:

```
SELECT DISTINCT select_list ...
```

(Instead of DISTINCT the key word ALL can be used to specify the default behavior of retaining all rows.)

Obviously, two rows are considered distinct if they differ in at least one column value. Null values are considered equal in this comparison.

Alternatively, an arbitrary expression can determine what rows are to be considered distinct:

```
SELECT DISTINCT ON (expression [, expression ...]) select_list ...
```

Here expression is an arbitrary value expression that is evaluated for all rows. A set of rows for which all the expressions are equal are considered duplicates, and only the first row of the set is kept in the output. Note that the “first row” of a set is unpredictable unless the query is sorted on enough columns to guarantee a unique ordering of the rows arriving at the DISTINCT filter. (DISTINCT ON processing occurs after ORDER BY sorting.)

The DISTINCT ON clause is not part of the SQL standard and is sometimes considered bad style because of the potentially indeterminate nature of its results. With judicious use of GROUP BY and subqueries in FROM, this construct can be avoided, but it is often the most convenient alternative.

### 7.4. Combining Queries

The results of two queries can be combined using the set operations union, intersection, and difference. The syntax is

```
query1 UNION [ALL] query2
query1 INTERSECT [ALL] query2
query1 EXCEPT [ALL] query2
```

query1 and query2 are queries that can use any of the features discussed up to this point. Set operations can also be nested and chained, for example

```
query1 UNION query2 UNION query3
```

which is executed as:

```
(query1 UNION query2) UNION query3
```
UNION effectively appends the result of query2 to the result of query1 (although there is no guarantee that this is the order in which the rows are actually returned). Furthermore, it eliminates duplicate rows from its result, in the same way as DISTINCT, unless UNION ALL is used.

INTERSECT returns all rows that are both in the result of query1 and in the result of query2. Duplicate rows are eliminated unless INTERSECT ALL is used.

EXCEPT returns all rows that are in the result of query1 but not in the result of query2. (This is sometimes called the difference between two queries.) Again, duplicates are eliminated unless EXCEPT ALL is used.

In order to calculate the union, intersection, or difference of two queries, the two queries must be “union compatible”, which means that they return the same number of columns and the corresponding columns have compatible data types, as described in Section 10.5.

7.5. Sorting Rows

After a query has produced an output table (after the select list has been processed) it can optionally be sorted. If sorting is not chosen, the rows will be returned in an unspecified order. The actual order in that case will depend on the scan and join plan types and the order on disk, but it must not be relied on. A particular output ordering can only be guaranteed if the sort step is explicitly chosen.

The ORDER BY clause specifies the sort order:

```
SELECT select_list
FROM table_expression
ORDER BY sort_expression1 [ASC | DESC] [NULLS { FIRST | LAST }]
[, sort_expression2 [ASC | DESC] [NULLS { FIRST | LAST }]] ...
```

The sort expression(s) can be any expression that would be valid in the query’s select list. An example is:

```
SELECT a, b FROM table1 ORDER BY a + b, c;
```

When more than one expression is specified, the later values are used to sort rows that are equal according to the earlier values. Each expression can be followed by an optional ASC or DESC keyword to set the sort direction to ascending or descending. ASC order is the default. Ascending order puts smaller values first, where “smaller” is defined in terms of the < operator. Similarly, descending order is determined with the > operator.¹

The NULLS FIRST and NULLS LAST options can be used to determine whether nulls appear before or after non-null values in the sort ordering. By default, null values sort as if larger than any non-null value; that is, NULLS FIRST is the default for DESC order, and NULLS LAST otherwise.

Note that the ordering options are considered independently for each sort column. For example ORDER BY x, y DESC means ORDER BY x ASC, y DESC, which is not the same as ORDER BY x DESC, y DESC.

A sort_expression can also be the column label or number of an output column, as in:

```
SELECT a + b AS sum, c FROM table1 ORDER BY sum;
SELECT a, max(b) FROM table1 GROUP BY a ORDER BY 1;
```

¹ Actually, PostgreSQL uses the default B-tree operator class for the expression’s data type to determine the sort ordering for ASC and DESC. Conventionally, data types will be set up so that the < and > operators correspond to this sort ordering, but a user-defined data type’s designer could choose to do something different.
both of which sort by the first output column. Note that an output column name has to stand alone, 
that is, it cannot be used in an expression — for example, this is not correct:

```
SELECT a + b AS sum, c FROM table1 ORDER BY sum + c; -- wrong
```

This restriction is made to reduce ambiguity. There is still ambiguity if an ORDER BY item is a simple 
name that could match either an output column name or a column from the table expression. The 
output column is used in such cases. This would only cause confusion if you use AS to rename an 
output column to match some other table column’s name.

ORDER BY can be applied to the result of a UNION, INTERSECT, or EXCEPT combination, but in this 
case it is only permitted to sort by output column names or numbers, not by expressions.

### 7.6. LIMIT and OFFSET

LIMIT and OFFSET allow you to retrieve just a portion of the rows that are generated by the rest of 
the query:

```
SELECT select_list
    FROM table_expression
    [ ORDER BY ... ]
    [ LIMIT { number | ALL } ] [ OFFSET number ]
```

If a limit count is given, no more than that many rows will be returned (but possibly fewer, if the query 
itself yields fewer rows). LIMIT ALL is the same as omitting the LIMIT clause, as is LIMIT with a 
NULL argument.

OFFSET says to skip that many rows before beginning to return rows. OFFSET 0 is the same as 
omitting the OFFSET clause, as is OFFSET with a NULL argument.

If both OFFSET and LIMIT appear, then OFFSET rows are skipped before starting to count the LIMIT 
rows that are returned.

When using LIMIT, it is important to use an ORDER BY clause that constrains the result rows into a 
unique order. Otherwise you will get an unpredictable subset of the query’s rows. You might be asking 
for the tenth through twentieth rows, but tenth through twentieth in what ordering? The ordering is 
unknown, unless you specified ORDER BY.

The query optimizer takes LIMIT into account when generating query plans, so you are very likely 
to get different plans (yielding different row orders) depending on what you give for LIMIT and 
OFFSET. Thus, using different LIMIT/OFFSET values to select different subsets of a query result will 
give inconsistent results unless you enforce a predictable result ordering with ORDER BY. This is not 
a bug; it is an inherent consequence of the fact that SQL does not promise to deliver the results of a 
query in any particular order unless ORDER BY is used to constrain the order.

The rows skipped by an OFFSET clause still have to be computed inside the server; therefore a large 
OFFSET might be inefficient.
7.7. VALUES Lists

VALUES provides a way to generate a “constant table” that can be used in a query without having to actually create and populate a table on-disk. The syntax is

```
VALUES ( expression [, ...] ) [ , ...]
```

Each parenthesized list of expressions generates a row in the table. The lists must all have the same number of elements (i.e., the number of columns in the table), and corresponding entries in each list must have compatible data types. The actual data type assigned to each column of the result is determined using the same rules as for UNION (see Section 10.5).

As an example:

```
VALUES (1, 'one'), (2, 'two'), (3, 'three');
```

will return a table of two columns and three rows. It’s effectively equivalent to:

```
SELECT 1 AS column1, 'one' AS column2
UNION ALL
SELECT 2, 'two'
UNION ALL
SELECT 3, 'three';
```

By default, PostgreSQL assigns the names `column1`, `column2`, etc. to the columns of a VALUES table. The column names are not specified by the SQL standard and different database systems do it differently, so it’s usually better to override the default names with a table alias list, like this:

```
SELECT * FROM (VALUES (1, 'one'), (2, 'two'), (3, 'three')) AS t (num, letter);
```

Syntactically, VALUES followed by expression lists is treated as equivalent to:

```
SELECT select_list FROM table_expression
```

and can appear anywhere a SELECT can. For example, you can use it as part of a UNION, or attach a `sort_specification` (ORDER BY, LIMIT, and/or OFFSET) to it. VALUES is most commonly used as the data source in an INSERT command, and next most commonly as a subquery.

For more information see VALUES.

7.8. WITH Queries (Common Table Expressions)

WITH provides a way to write auxiliary statements for use in a larger query. These statements, which are often referred to as Common Table Expressions or CTEs, can be thought of as defining temporary tables that exist just for one query. Each auxiliary statement in a WITH clause can be a SELECT, INSERT, UPDATE, or DELETE; and the WITH clause itself is attached to a primary statement that can also be a SELECT, INSERT, UPDATE, or DELETE.
Chapter 7. Queries

7.8.1. SELECT in WITH

The basic value of SELECT in WITH is to break down complicated queries into simpler parts. An example is:

```sql
WITH regional_sales AS (
    SELECT region, SUM(amount) AS total_sales
    FROM orders
    GROUP BY region
), top_regions AS (
    SELECT region
    FROM regional_sales
    WHERE total_sales > (SELECT SUM(total_sales)/10 FROM regional_sales)
)
SELECT region,
    product,
    SUM(quantity) AS product_units,
    SUM(amount) AS product_sales
FROM orders
WHERE region IN (SELECT region FROM top_regions)
GROUP BY region, product;
```

which displays per-product sales totals in only the top sales regions. The WITH clause defines two auxiliary statements named `regional_sales` and `top_regions`, where the output of `regional_sales` is used in `top_regions` and the output of `top_regions` is used in the primary `SELECT` query. This example could have been written without WITH, but we’d have needed two levels of nested sub-SELECTs. It’s a bit easier to follow this way.

The optional RECURSIVE modifier changes WITH from a mere syntactic convenience into a feature that accomplishes things not otherwise possible in standard SQL. Using RECURSIVE, a WITH query can refer to its own output. A very simple example is this query to sum the integers from 1 through 100:

```sql
WITH RECURSIVE t(n) AS (
    VALUES (1)
    UNION ALL
    SELECT n+1 FROM t WHERE n < 100
)
SELECT sum(n) FROM t;
```

The general form of a recursive WITH query is always a non-recursive term, then UNION (or UNION ALL), then a recursive term, where only the recursive term can contain a reference to the query’s own output. Such a query is executed as follows:

Recursive Query Evaluation

1. Evaluate the non-recursive term. For UNION (but not UNION ALL), discard duplicate rows. Include all remaining rows in the result of the recursive query, and also place them in a temporary working table.
2. So long as the working table is not empty, repeat these steps:
   a. Evaluate the recursive term, substituting the current contents of the working table for the recursive self-reference. For UNION (but not UNION ALL), discard duplicate rows and rows that duplicate any previous result row. Include all remaining rows in the result of the recursive query, and also place them in a temporary intermediate table.
Chapter 7. Queries

b. Replace the contents of the working table with the contents of the intermediate table, then empty the intermediate table.

**Note:** Strictly speaking, this process is iteration not recursion, but **RECURSIVE** is the terminology chosen by the SQL standards committee.

In the example above, the working table has just a single row in each step, and it takes on the values from 1 through 100 in successive steps. In the 100th step, there is no output because of the **WHERE** clause, and so the query terminates.

Recursive queries are typically used to deal with hierarchical or tree-structured data. A useful example is this query to find all the direct and indirect sub-parts of a product, given only a table that shows immediate inclusions:

```
WITH RECURSIVE included_parts(sub_part, part, quantity) AS (
  SELECT sub_part, part, quantity FROM parts WHERE part = 'our_product'
  UNION ALL
  SELECT p.sub_part, p.part, p.quantity
  FROM included_parts pr, parts p
  WHERE p.part = pr.sub_part
)
SELECT sub_part, SUM(quantity) as total_quantity
FROM included_parts
GROUP BY sub_part
```

When working with recursive queries it is important to be sure that the recursive part of the query will eventually return no tuples, or else the query will loop indefinitely. Sometimes, using **UNION** instead of **UNION ALL** can accomplish this by discarding rows that duplicate previous output rows. However, often a cycle does not involve output rows that are completely duplicate: it may be necessary to check just one or a few fields to see if the same point has been reached before. The standard method for handling such situations is to compute an array of the already-visited values. For example, consider the following query that searches a table **graph** using a **link** field:

```
WITH RECURSIVE search_graph(id, link, data, depth) AS (
  SELECT g.id, g.link, g.data, 1
  FROM graph g
  UNION ALL
  SELECT g.id, g.link, g.data, sg.depth + 1
  FROM graph g, search_graph sg
  WHERE g.id = sg.link
)
SELECT * FROM search_graph;
```

This query will loop if the **link** relationships contain cycles. Because we require a “depth” output, just changing **UNION ALL** to **UNION** would not eliminate the looping. Instead we need to recognize whether we have reached the same row again while following a particular path of links. We add two columns **path** and **cycle** to the loop-prone query:

```
WITH RECURSIVE search_graph(id, link, data, depth, path, cycle) AS (
  SELECT g.id, g.link, g.data, 1,
  ARRAY[g.id],
  false
  UNION ALL
  SELECT g.id, g.link, g.data, sg.depth + 1,
  SUB_ARRAY(path, 1, sg.depth + 1) || g.id,
  false
  FROM graph g, search_graph sg
  WHERE g.id = sg.link
)
SELECT * FROM search_graph;
```

114
Chapter 7. Queries

```sql
FROM graph g
UNION ALL
SELECT g.id, g.link, g.data, sg.depth + 1,
    path || g.id,
    g.id = ANY(path)
FROM graph g, search_graph sg
WHERE g.id = sg.link AND NOT cycle
)
SELECT * FROM search_graph;

Aside from preventing cycles, the array value is often useful in its own right as representing the “path” taken to reach any particular row.

In the general case where more than one field needs to be checked to recognize a cycle, use an array of rows. For example, if we needed to compare fields f1 and f2:

```sql
WITH RECURSIVE search_graph(id, link, data, depth, path, cycle) AS (  
    SELECT g.id, g.link, g.data, 1,
        ARRAY[ROW(g.f1, g.f2)],
        false
    FROM graph g
    UNION ALL
    SELECT g.id, g.link, g.data, sg.depth + 1,
        path || ROW(g.f1, g.f2),
        ROW(g.f1, g.f2) = ANY(path)
    FROM graph g, search_graph sg
    WHERE g.id = sg.link AND NOT cycle
)
SELECT * FROM search_graph;
```

**Tip:** Omit the `ROW()` syntax in the common case where only one field needs to be checked to recognize a cycle. This allows a simple array rather than a composite-type array to be used, gaining efficiency.

**Tip:** The recursive query evaluation algorithm produces its output in breadth-first search order. You can display the results in depth-first search order by making the outer query `ORDER BY` a “path” column constructed in this way.

A helpful trick for testing queries when you are not certain if they might loop is to place a `LIMIT` in the parent query. For example, this query would loop forever without the `LIMIT`:

```sql
WITH RECURSIVE t(n) AS (  
    SELECT 1
    UNION ALL
    SELECT n+1 FROM t
)
SELECT n FROM t LIMIT 100;
```

This works because PostgreSQL’s implementation evaluates only as many rows of a `WITH` query as are actually fetched by the parent query. Using this trick in production is not recommended, because other systems might work differently. Also, it usually won’t work if you make the outer query sort the
recursive query’s results or join them to some other table, because in such cases the outer query will usually try to fetch all of the WITH query’s output anyway.

A useful property of WITH queries is that they are evaluated only once per execution of the parent query, even if they are referred to more than once by the parent query or sibling WITH queries. Thus, expensive calculations that are needed in multiple places can be placed within a WITH query to avoid redundant work. Another possible application is to prevent unwanted multiple evaluations of functions with side-effects. However, the other side of this coin is that the optimizer is less able to push restrictions from the parent query down into a WITH query than an ordinary subquery. The WITH query will generally be evaluated as written, without suppression of rows that the parent query might discard afterwards. (But, as mentioned above, evaluation might stop early if the reference(s) to the query demand only a limited number of rows.)

The examples above only show WITH being used with SELECT, but it can be attached in the same way to INSERT, UPDATE, or DELETE. In each case it effectively provides temporary table(s) that can be referred to in the main command.

### 7.8.2. Data-Modifying Statements in WITH

You can use data-modifying statements (INSERT, UPDATE, or DELETE) in WITH. This allows you to perform several different operations in the same query. An example is:

```sql
WITH moved_rows AS (
    DELETE FROM products
    WHERE
        "date" >= '2010-10-01' AND
        "date" < '2010-11-01'
    RETURNING *
)
INSERT INTO products_log
SELECT * FROM moved_rows;
```

This query effectively moves rows from products to products_log. The DELETE in WITH deletes the specified rows from products, returning their contents by means of its RETURNING clause; and then the primary query reads that output and inserts it into products_log.

A fine point of the above example is that the WITH clause is attached to the INSERT, not the sub-SELECT within the INSERT. This is necessary because data-modifying statements are only allowed in WITH clauses that are attached to the top-level statement. However, normal WITH visibility rules apply, so it is possible to refer to the WITH statement’s output from the sub-SELECT.

Data-modifying statements in WITH usually have RETURNING clauses (see Section 6.4), as shown in the example above. It is the output of the RETURNING clause, not the target table of the data-modifying statement, that forms the temporary table that can be referred to by the rest of the query. If a data-modifying statement in WITH lacks a RETURNING clause, then it forms no temporary table and cannot be referred to in the rest of the query. Such a statement will be executed nonetheless. A not-particularly-useful example is:

```sql
WITH t AS (  
    DELETE FROM foo
)
DELETE FROM bar;
```

This example would remove all rows from tables foo and bar. The number of affected rows reported to the client would only include rows removed from bar.
Recursive self-references in data-modifying statements are not allowed. In some cases it is possible to work around this limitation by referring to the output of a recursive WITH, for example:

```
WITH RECURSIVE included_parts(sub_part, part) AS (  
    SELECT sub_part, part FROM parts WHERE part = 'our_product'  
    UNION ALL  
    SELECT p.sub_part, p.part  
    FROM included_parts pr, parts p  
    WHERE p.part = pr.sub_part
)
DELETE FROM parts  
WHERE part IN (SELECT part FROM included_parts);
```

This query would remove all direct and indirect subparts of a product.

Data-modifying statements in WITH are executed exactly once, and always to completion, independently of whether the primary query reads all (or indeed any) of their output. Notice that this is different from the rule for SELECT in WITH: as stated in the previous section, execution of a SELECT is carried only as far as the primary query demands its output.

The sub-statements in WITH are executed concurrently with each other and with the main query. Therefore, when using data-modifying statements in WITH, the order in which the specified updates actually happen is unpredictable. All the statements are executed with the same snapshot (see Chapter 13), so they cannot “see” one another’s effects on the target tables. This alleviates the effects of the unpredictability of the actual order of row updates, and means that RETURNING data is the only way to communicate changes between different WITH sub-statements and the main query. An example of this is that in

```
WITH t AS (  
    UPDATE products SET price = price * 1.05  
    RETURNING *
)
SELECT * FROM products;
```

the outer SELECT would return the original prices before the action of the UPDATE, while in

```
WITH t AS (  
    UPDATE products SET price = price * 1.05  
    RETURNING *
)
SELECT * FROM t;
```

the outer SELECT would return the updated data.

Trying to update the same row twice in a single statement is not supported. Only one of the modifications takes place, but it is not easy (and sometimes not possible) to reliably predict which one. This also applies to deleting a row that was already updated in the same statement: only the update is performed. Therefore you should generally avoid trying to modify a single row twice in a single statement. In particular avoid writing WITH sub-statements that could affect the same rows changed by the main statement or a sibling sub-statement. The effects of such a statement will not be predictable.

At present, any table used as the target of a data-modifying statement in WITH must not have a conditional rule, nor an ALSO rule, nor an INSTEAD rule that expands to multiple statements.
Chapter 8. Data Types

PostgreSQL has a rich set of native data types available to users. Users can add new types to PostgreSQL using the CREATE TYPE command.

Table 8-1 shows all the built-in general-purpose data types. Most of the alternative names listed in the “Aliases” column are the names used internally by PostgreSQL for historical reasons. In addition, some internally used or deprecated types are available, but are not listed here.

<table>
<thead>
<tr>
<th>Name</th>
<th>Aliases</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bigint</td>
<td>int8</td>
<td>signed eight-byte integer</td>
</tr>
<tr>
<td>bigserial</td>
<td>serial8</td>
<td>autoincrementing eight-byte integer</td>
</tr>
<tr>
<td>bit [ (n) ]</td>
<td></td>
<td>fixed-length bit string</td>
</tr>
<tr>
<td>bit varying [ (n) ]</td>
<td>varbit [ (n) ]</td>
<td>variable-length bit string</td>
</tr>
<tr>
<td>boolean</td>
<td>bool</td>
<td>logical Boolean (true/false)</td>
</tr>
<tr>
<td>box</td>
<td></td>
<td>rectangular box on a plane</td>
</tr>
<tr>
<td>bytea</td>
<td></td>
<td>binary data (“byte array”)</td>
</tr>
<tr>
<td>character [ (n) ]</td>
<td>char [ (n) ]</td>
<td>fixed-length character string</td>
</tr>
<tr>
<td>character varying [ (n) ]</td>
<td>varchar [ (n) ]</td>
<td>variable-length character string</td>
</tr>
<tr>
<td>cidr</td>
<td></td>
<td>IPv4 or IPv6 network address</td>
</tr>
<tr>
<td>circle</td>
<td></td>
<td>circle on a plane</td>
</tr>
<tr>
<td>date</td>
<td></td>
<td>calendar date (year, month, day)</td>
</tr>
<tr>
<td>double precision</td>
<td>float8</td>
<td>double precision floating-point number (8 bytes)</td>
</tr>
<tr>
<td>inet</td>
<td></td>
<td>IPv4 or IPv6 host address</td>
</tr>
<tr>
<td>integer</td>
<td>int, int4</td>
<td>signed four-byte integer</td>
</tr>
<tr>
<td>interval [ fields ] [ (p) ]</td>
<td></td>
<td>time span</td>
</tr>
<tr>
<td>json</td>
<td></td>
<td>textual JSON data</td>
</tr>
<tr>
<td>jsonb</td>
<td></td>
<td>binary JSON data, decomposed</td>
</tr>
<tr>
<td>line</td>
<td></td>
<td>infinite line on a plane</td>
</tr>
<tr>
<td>lseg</td>
<td></td>
<td>line segment on a plane</td>
</tr>
<tr>
<td>macaddr</td>
<td></td>
<td>MAC (Media Access Control) address</td>
</tr>
<tr>
<td>money</td>
<td></td>
<td>currency amount</td>
</tr>
<tr>
<td>numeric [ (p, s) ]</td>
<td>decimal [ (p, s) ]</td>
<td>exact numeric of selectable precision</td>
</tr>
<tr>
<td>path</td>
<td></td>
<td>geometric path on a plane</td>
</tr>
<tr>
<td>pg_lsn</td>
<td></td>
<td>PostgreSQL Log Sequence Number</td>
</tr>
<tr>
<td>point</td>
<td></td>
<td>geometric point on a plane</td>
</tr>
</tbody>
</table>
Chapter 8. Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Aliases</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>polygon</td>
<td></td>
<td>closed geometric path on a plane</td>
</tr>
<tr>
<td>real</td>
<td>float4</td>
<td>single precision floating-point number (4 bytes)</td>
</tr>
<tr>
<td>smallint</td>
<td>int2</td>
<td>signed two-byte integer</td>
</tr>
<tr>
<td>smallserial</td>
<td></td>
<td>autoincrementing two-byte integer</td>
</tr>
<tr>
<td>serial</td>
<td>serial4</td>
<td>autoincrementing four-byte integer</td>
</tr>
<tr>
<td>text</td>
<td></td>
<td>variable-length character string</td>
</tr>
<tr>
<td>time [ (p) ] [ without time zone ]</td>
<td></td>
<td>time of day (no time zone)</td>
</tr>
<tr>
<td>time [ (p) ] with time zone</td>
<td>timetz</td>
<td>time of day, including time zone</td>
</tr>
<tr>
<td>timestamp [ (p) ] [ without time zone ]</td>
<td></td>
<td>date and time (no time zone)</td>
</tr>
<tr>
<td>timestamp [ (p) ] with time zone</td>
<td>timestamptz</td>
<td>date and time, including time zone</td>
</tr>
<tr>
<td>tsquery</td>
<td></td>
<td>text search query</td>
</tr>
<tr>
<td>tsvector</td>
<td></td>
<td>text search document</td>
</tr>
<tr>
<td>txid_snapshot</td>
<td></td>
<td>user-level transaction ID snapshot</td>
</tr>
<tr>
<td>uuid</td>
<td></td>
<td>universally unique identifier</td>
</tr>
<tr>
<td>xml</td>
<td></td>
<td>XML data</td>
</tr>
</tbody>
</table>

Compatibility: The following types (or spellings thereof) are specified by SQL: bigint, bit, bit varying, boolean, char, character varying, character, varchar, date, double precision, integer, interval, numeric, decimal, real, smallint, time (with or without time zone), timestamp (with or without time zone), xml.

Each data type has an external representation determined by its input and output functions. Many of the built-in types have obvious external formats. However, several types are either unique to PostgreSQL, such as geometric paths, or have several possible formats, such as the date and time types. Some of the input and output functions are not invertible, i.e., the result of an output function might lose accuracy when compared to the original input.

8.1. Numeric Types

Numeric types consist of two-, four-, and eight-byte integers, four- and eight-byte floating-point numbers, and selectable-precision decimals. Table 8-2 lists the available types.

Table 8-2. Numeric Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>polygon</td>
<td></td>
<td>closed geometric path on a plane</td>
</tr>
<tr>
<td>real</td>
<td>float4</td>
<td>single precision floating-point number (4 bytes)</td>
</tr>
<tr>
<td>smallint</td>
<td>int2</td>
<td>signed two-byte integer</td>
</tr>
<tr>
<td>smallserial</td>
<td>serial2</td>
<td>autoincrementing two-byte integer</td>
</tr>
<tr>
<td>serial</td>
<td>serial4</td>
<td>autoincrementing four-byte integer</td>
</tr>
<tr>
<td>text</td>
<td></td>
<td>variable-length character string</td>
</tr>
<tr>
<td>time</td>
<td></td>
<td>time of day (no time zone)</td>
</tr>
<tr>
<td>time with</td>
<td>timetz</td>
<td>time of day, including time zone</td>
</tr>
<tr>
<td>timestamp</td>
<td>timestamptz</td>
<td>date and time, including time zone</td>
</tr>
<tr>
<td>tsquery</td>
<td></td>
<td>text search query</td>
</tr>
<tr>
<td>tsvector</td>
<td></td>
<td>text search document</td>
</tr>
<tr>
<td>txid_snapshot</td>
<td></td>
<td>user-level transaction ID snapshot</td>
</tr>
<tr>
<td>uuid</td>
<td></td>
<td>universally unique identifier</td>
</tr>
<tr>
<td>xml</td>
<td></td>
<td>XML data</td>
</tr>
</tbody>
</table>
Chapter 8. Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Range</th>
</tr>
</thead>
<tbody>
<tr>
<td>smallint</td>
<td>2 bytes</td>
<td>small-range integer</td>
<td>-32768 to +32767</td>
</tr>
<tr>
<td>integer</td>
<td>4 bytes</td>
<td>typical choice for integer</td>
<td>-2147483648 to +2147483647</td>
</tr>
<tr>
<td>bigint</td>
<td>8 bytes</td>
<td>large-range integer</td>
<td>-9223372036854775808 to +9223372036854775807</td>
</tr>
<tr>
<td>decimal</td>
<td>variable</td>
<td>user-specified precision, exact</td>
<td>up to 131072 digits before the decimal point; up to 16383 digits after the decimal point</td>
</tr>
<tr>
<td>numeric</td>
<td>variable</td>
<td>user-specified precision, exact</td>
<td>up to 131072 digits before the decimal point; up to 16383 digits after the decimal point</td>
</tr>
<tr>
<td>real</td>
<td>4 bytes</td>
<td>variable-precision, inexact</td>
<td>6 decimal digits precision</td>
</tr>
<tr>
<td>double precision</td>
<td>8 bytes</td>
<td>variable-precision, inexact</td>
<td>15 decimal digits precision</td>
</tr>
<tr>
<td>smallserial</td>
<td>2 bytes</td>
<td>small autoincrementing integer</td>
<td>1 to 32767</td>
</tr>
<tr>
<td>serial</td>
<td>4 bytes</td>
<td>autoincrementing integer</td>
<td>1 to 2147483647</td>
</tr>
<tr>
<td>bigserial</td>
<td>8 bytes</td>
<td>large autoincrementing integer</td>
<td>1 to 9223372036854775807</td>
</tr>
</tbody>
</table>

The syntax of constants for the numeric types is described in Section 4.1.2. The numeric types have a full set of corresponding arithmetic operators and functions. Refer to Chapter 9 for more information. The following sections describe the types in detail.

### 8.1.1. Integer Types

The types smallint, integer, and bigint store whole numbers, that is, numbers without fractional components, of various ranges. Attempts to store values outside of the allowed range will result in an error.

The type integer is the common choice, as it offers the best balance between range, storage size, and performance. The smallint type is generally only used if disk space is at a premium. The bigint type is designed to be used when the range of the integer type is insufficient.

SQL only specifies the integer types integer (or int), smallint, and bigint. The type names int2, int4, and int8 are extensions, which are also used by some other SQL database systems.
8.1.2. Arbitrary Precision Numbers

The type `numeric` can store numbers with a very large number of digits. It is especially recommended for storing monetary amounts and other quantities where exactness is required. Calculations with `numeric` values yield exact results where possible, e.g., addition, subtraction, multiplication. However, calculations on `numeric` values are very slow compared to the integer types, or to the floating-point types described in the next section.

We use the following terms below: the **precision** of a `numeric` is the total count of significant digits in the whole number, that is, the number of digits to both sides of the decimal point. The **scale** of a `numeric` is the count of decimal digits in the fractional part, to the right of the decimal point. So the number 23.5141 has a precision of 6 and a scale of 4. Integers can be considered to have a scale of zero.

Both the maximum precision and the maximum scale of a `numeric` column can be configured. To declare a column of type `numeric` use the syntax:

```
NUMERIC(precision, scale)
```

The precision must be positive, the scale zero or positive. Alternatively:

```
NUMERIC(precision)
```

selects a scale of 0. Specifying:

```
NUMERIC
```

without any precision or scale creates a column in which numeric values of any precision and scale can be stored, up to the implementation limit on precision. A column of this kind will not coerce input values to any particular scale, whereas `numeric` columns with a declared scale will coerce input values to that scale. (The SQL standard requires a default scale of 0, i.e., coercion to integer precision. We find this a bit useless. If you’re concerned about portability, always specify the precision and scale explicitly.)

**Note:** The maximum allowed precision when explicitly specified in the type declaration is 1000; `NUMERIC` without a specified precision is subject to the limits described in Table 8-2.

If the scale of a value to be stored is greater than the declared scale of the column, the system will round the value to the specified number of fractional digits. Then, if the number of digits to the left of the decimal point exceeds the declared precision minus the declared scale, an error is raised.

Numeric values are physically stored without any extra leading or trailing zeroes. Thus, the declared precision and scale of a column are maximums, not fixed allocations. (In this sense the `numeric` type is more akin to `varchar(n)` than to `char(n)`.) The actual storage requirement is two bytes for each group of four decimal digits, plus three to eight bytes overhead.

In addition to ordinary numeric values, the `numeric` type allows the special value `NaN`, meaning “not-a-number”. Any operation on `NaN` yields another `NaN`. When writing this value as a constant in an SQL command, you must put quotes around it, for example `UPDATE table SET x = 'NaN'`. On input, the string `NaN` is recognized in a case-insensitive manner.

**Note:** In most implementations of the “not-a-number” concept, `NaN` is not considered equal to any other numeric value (including `NaN`). In order to allow `numeric` values to be sorted and used in tree-based indexes, PostgreSQL treats `NaN` values as equal, and greater than all non-`NaN` values.
The types \texttt{decimal} and \texttt{numeric} are equivalent. Both types are part of the SQL standard.

When rounding values, the \texttt{numeric} type rounds ties away from zero, while (on most machines) the \texttt{real} and \texttt{double precision} types round ties to the nearest even number. For example:

\begin{verbatim}
SELECT x,
       round(x::numeric) AS num_round,
       round(x::double precision) AS dbl_round
FROM generate_series(-3.5, 3.5, 1) as x;
\end{verbatim}

\begin{tabular}{c|c|c}
\textbf{x} & \textbf{num_round} & \textbf{dbl_round} \\
\hline
-3.5 & -4 & -4 \\
-2.5 & -3 & -2 \\
-1.5 & -2 & -2 \\
-0.5 & -1 & -0 \\
0.5 & 1 & 0 \\
1.5 & 2 & 2 \\
2.5 & 3 & 2 \\
3.5 & 4 & 4 \\
\hline
\end{tabular}

(8 rows)

\subsection{8.1.3. Floating-Point Types}

The data types \texttt{real} and \texttt{double precision} are inexact, variable-precision numeric types. In practice, these types are usually implementations of IEEE Standard 754 for Binary Floating-Point Arithmetic (single and double precision, respectively), to the extent that the underlying processor, operating system, and compiler support it.

Inexact means that some values cannot be converted exactly to the internal format and are stored as approximations, so that storing and retrieving a value might show slight discrepancies. Managing these errors and how they propagate through calculations is the subject of an entire branch of mathematics and computer science and will not be discussed here, except for the following points:

\begin{itemize}
\item If you require exact storage and calculations (such as for monetary amounts), use the \texttt{numeric} type instead.
\item If you want to do complicated calculations with these types for anything important, especially if you rely on certain behavior in boundary cases (infinity, underflow), you should evaluate the implementation carefully.
\item Comparing two floating-point values for equality might not always work as expected.
\end{itemize}

On most platforms, the \texttt{real} type has a range of at least 1E-37 to 1E+37 with a precision of at least 6 decimal digits. The \texttt{double precision} type typically has a range of around 1E-307 to 1E+308 with a precision of at least 15 digits. Values that are too large or too small will cause an error. Rounding might take place if the precision of an input number is too high. Numbers too close to zero that are not representable as distinct from zero will cause an underflow error.

\textbf{Note}: The \texttt{extra_float_digits} setting controls the number of extra significant digits included when a floating point value is converted to text for output. With the default value of 0, the output is the same on every platform supported by PostgreSQL. Increasing it will produce output that more accurately represents the stored value, but may be unportable.
In addition to ordinary numeric values, the floating-point types have several special values:

- **Infinity**
- **-Infinity**
- **NaN**

These represent the IEEE 754 special values “infinity”, “negative infinity”, and “not-a-number”, respectively. (On a machine whose floating-point arithmetic does not follow IEEE 754, these values will probably not work as expected.) When writing these values as constants in an SQL command, you must put quotes around them, for example `UPDATE table SET x = 'Infinity'`. On input, these strings are recognized in a case-insensitive manner.

**Note:** IEEE754 specifies that NaN should not compare equal to any other floating-point value (including NaN). In order to allow floating-point values to be sorted and used in tree-based indexes, PostgreSQL treats NaN values as equal, and greater than all non-NaN values.

PostgreSQL also supports the SQL-standard notations `float` and `float(p)` for specifying inexact numeric types. Here, `p` specifies the minimum acceptable precision in binary digits. PostgreSQL accepts `float(1)` to `float(24)` as selecting the real type, while `float(25)` to `float(53)` select double precision. Values of `p` outside the allowed range draw an error. `float` with no precision specified is taken to mean double precision.

**Note:** The assumption that real and double precision have exactly 24 and 53 bits in the mantissa respectively is correct for IEEE-standard floating point implementations. On non-IEEE platforms it might be off a little, but for simplicity the same ranges of `p` are used on all platforms.

### 8.1.4. Serial Types

The data types `smallserial`, `serial` and `bigserial` are not true types, but merely a notational convenience for creating unique identifier columns (similar to the AUTO_INCREMENT property supported by some other databases). In the current implementation, specifying:

```sql
CREATE TABLE tablename (
  colname SERIAL
);
```

is equivalent to specifying:

```sql
CREATE SEQUENCE tablename_colname_seq;
CREATE TABLE tablename (
  colname integer NOT NULL DEFAULT nextval('tablename_colname_seq')
);
ALTER SEQUENCE tablename_colname_seq OWNED BY tablename.colname;
```

Thus, we have created an integer column and arranged for its default values to be assigned from a sequence generator. A `NOT NULL` constraint is applied to ensure that a null value cannot be inserted. (In most cases you would also want to attach a `UNIQUE` or `PRIMARY KEY` constraint to prevent duplicate
values from being inserted by accident, but this is not automatic.) Lastly, the sequence is marked as “owned by” the column, so that it will be dropped if the column or table is dropped.

**Note:** Because `smallserial`, `serial` and `bigserial` are implemented using sequences, there may be “holes” or gaps in the sequence of values which appears in the column, even if no rows are ever deleted. A value allocated from the sequence is still “used up” even if a row containing that value is never successfully inserted into the table column. This may happen, for example, if the inserting transaction rolls back. See `nextval()` in Section 9.16 for details.

To insert the next value of the sequence into the `serial` column, specify that the `serial` column should be assigned its default value. This can be done either by excluding the column from the list of columns in the `INSERT` statement, or through the use of the `DEFAULT` key word.

The type names `serial` and `serial4` are equivalent: both create integer columns. The type names `bigserial` and `serial8` work the same way, except that they create a bigint column. `bigserial` should be used if you anticipate the use of more than $2^{31}$ identifiers over the lifetime of the table. The type names `smallserial` and `serial2` also work the same way, except that they create a smallint column.

The sequence created for a `serial` column is automatically dropped when the owning column is dropped. You can drop the sequence without dropping the column, but this will force removal of the column default expression.

## 8.2. Monetary Types

The `money` type stores a currency amount with a fixed fractional precision; see Table 8-3. The fractional precision is determined by the database’s `lc_monetary` setting. The range shown in the table assumes there are two fractional digits. Input is accepted in a variety of formats, including integer and floating-point literals, as well as typical currency formatting, such as ‘$1,000.00’. Output is generally in the latter form but depends on the locale.

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Range</th>
</tr>
</thead>
<tbody>
<tr>
<td>money</td>
<td>8 bytes</td>
<td>currency amount</td>
<td>-92233720368547758.08 to +92233720368547758.07</td>
</tr>
</tbody>
</table>

Since the output of this data type is locale-sensitive, it might not work to load `money` data into a database that has a different setting of `lc_monetary`. To avoid problems, before restoring a dump into a new database make sure `lc_monetary` has the same or equivalent value as in the database that was dumped.

Values of the `numeric`, `int`, and `bigint` data types can be cast to `money`. Conversion from the `real` and `double precision` data types can be done by casting to `numeric` first, for example:

```sql
SELECT '12.34'::float8::numeric::money;
```

However, this is not recommended. Floating point numbers should not be used to handle money due
to the potential for rounding errors.

A money value can be cast to numeric without loss of precision. Conversion to other types could potentially lose precision, and must also be done in two stages:

```
SELECT '52093.89'::money::numeric::float8;
```

Division of a money value by an integer value is performed with truncation of the fractional part towards zero. To get a rounded result, divide by a floating-point value, or cast the money value to numeric before dividing and back to money afterwards. (The latter is preferable to avoid risking precision loss.) When a money value is divided by another money value, the result is double precision (i.e., a pure number, not money); the currency units cancel each other out in the division.

## 8.3. Character Types

Table 8-4. Character Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>character varying(n), varchar(n)</td>
<td>variable-length with limit</td>
</tr>
<tr>
<td>character(n), char(n)</td>
<td>fixed-length, blank padded</td>
</tr>
<tr>
<td>text</td>
<td>variable unlimited length</td>
</tr>
</tbody>
</table>

Table 8-4 shows the general-purpose character types available in PostgreSQL.

SQL defines two primary character types: character varying(n) and character(n), where n is a positive integer. Both of these types can store strings up to n characters (not bytes) in length. An attempt to store a longer string into a column of these types will result in an error, unless the excess characters are all spaces, in which case the string will be truncated to the maximum length. (This somewhat bizarre exception is required by the SQL standard.) If the string to be stored is shorter than the declared length, values of type character will be space-padded; values of type character varying will simply store the shorter string.

If one explicitly casts a value to character varying(n) or character(n), then an over-length value will be truncated to n characters without raising an error. (This too is required by the SQL standard.)

The notations varchar(n) and char(n) are aliases for character varying(n) and character(n), respectively. character without length specifier is equivalent to character(1). If character varying is used without length specifier, the type accepts strings of any size. The latter is a PostgreSQL extension.

In addition, PostgreSQL provides the text type, which stores strings of any length. Although the type text is not in the SQL standard, several other SQL database management systems have it as well.

Values of type character are physically padded with spaces to the specified width n, and are stored and displayed that way. However, trailing spaces are treated as semantically insignificant and disregarded when comparing two values of type character. In collations where whitespace is significant, this behavior can produce unexpected results; for example SELECT 'a'::CHAR(2) collate "C" < E'a\n'::CHAR(2) returns true, even though C locale would consider a space to be greater than a newline. Trailing spaces are removed when converting a character value to one of the other string types. Note that trailing spaces are semantically significant in character varying and text val-
ues, and when using pattern matching, that is LIKE and regular expressions.

The storage requirement for a short string (up to 126 bytes) is 1 byte plus the actual string, which includes the space padding in the case of character. Longer strings have 4 bytes of overhead instead of 1. Long strings are compressed by the system automatically, so the physical requirement on disk might be less. Very long values are also stored in background tables so that they do not interfere with rapid access to shorter column values. In any case, the longest possible character string that can be stored is about 1 GB. (The maximum value that will be allowed for \( n \) in the data type declaration is less than that. It wouldn’t be useful to change this because with multibyte character encodings the number of characters and bytes can be quite different. If you desire to store long strings with no specific upper limit, use text or character varying without a length specifier, rather than making up an arbitrary length limit.)

**Tip:** There is no performance difference among these three types, apart from increased storage space when using the blank-padded type, and a few extra CPU cycles to check the length when storing into a length-constrained column. While character(\( n \)) has performance advantages in some other database systems, there is no such advantage in PostgreSQL; in fact character(\( n \)) is usually the slowest of the three because of its additional storage costs. In most situations text or character varying should be used instead.

Refer to Section 4.1.2.1 for information about the syntax of string literals, and to Chapter 9 for information about available operators and functions. The database character set determines the character set used to store textual values; for more information on character set support, refer to Section 23.3.

### Example 8-1. Using the Character Types

```sql
CREATE TABLE test1 (a character(4));
INSERT INTO test1 VALUES ('ok');
SELECT a, char_length(a) FROM test1; -- ①

<table>
<thead>
<tr>
<th>a</th>
<th>char_length</th>
</tr>
</thead>
<tbody>
<tr>
<td>ok</td>
<td>2</td>
</tr>
</tbody>
</table>
```

```sql
CREATE TABLE test2 (b varchar(5));
INSERT INTO test2 VALUES ('ok');
INSERT INTO test2 VALUES ('good');
INSERT INTO test2 VALUES ('too long');
ERROR: value too long for type character varying(5)
INSERT INTO test2 VALUES ('too long':varchar(5)); -- explicit truncation
SELECT b, char_length(b) FROM test2;

<table>
<thead>
<tr>
<th>b</th>
<th>char_length</th>
</tr>
</thead>
<tbody>
<tr>
<td>ok</td>
<td>2</td>
</tr>
<tr>
<td>good</td>
<td>5</td>
</tr>
<tr>
<td>too l</td>
<td>5</td>
</tr>
</tbody>
</table>
```

① The `char_length` function is discussed in Section 9.4.

There are two other fixed-length character types in PostgreSQL, shown in Table 8-5. The `name` type exists only for the storage of identifiers in the internal system catalogs and is not intended for use by the general user. Its length is currently defined as 64 bytes (63 usable characters plus terminator) but should be referenced using the constant `NAMEDATALEN` in C source code. The length is set at compile time (and is therefore adjustable for special uses); the default maximum length might change in a
future release. The type "char" (note the quotes) is different from char(1) in that it only uses one byte of storage. It is internally used in the system catalogs as a simplistic enumeration type.

<p>| Table 8-5. Special Character Types |</p>
<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>&quot;char&quot;</td>
<td>1 byte</td>
<td>single-byte internal type</td>
</tr>
<tr>
<td>name</td>
<td>64 bytes</td>
<td>internal type for object names</td>
</tr>
</tbody>
</table>

### 8.4. Binary Data Types

The bytea data type allows storage of binary strings; see Table 8-6.

<p>| Table 8-6. Binary Data Types |</p>
<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bytea</td>
<td>1 or 4 bytes plus the actual binary string</td>
<td>variable-length binary string</td>
</tr>
</tbody>
</table>

A binary string is a sequence of octets (or bytes). Binary strings are distinguished from character strings in two ways. First, binary strings specifically allow storing octets of value zero and other “non-printable” octets (usually, octets outside the decimal range 32 to 126). Character strings disallow zero octets, and also disallow any other octet values and sequences of octet values that are invalid according to the database’s selected character set encoding. Second, operations on binary strings process the actual bytes, whereas the processing of character strings depends on locale settings. In short, binary strings are appropriate for storing data that the programmer thinks of as “raw bytes”, whereas character strings are appropriate for storing text.

The bytea type supports two formats for input and output: “hex” format and PostgreSQL’s historical “escape” format. Both of these are always accepted on input. The output format depends on the configuration parameter bytea_output; the default is hex. (Note that the hex format was introduced in PostgreSQL 9.0; earlier versions and some tools don’t understand it.)

The SQL standard defines a different binary string type, called BLOB or BINARY LARGE OBJECT. The input format is different from bytea, but the provided functions and operators are mostly the same.

#### 8.4.1. bytea Hex Format

The “hex” format encodes binary data as 2 hexadecimal digits per byte, most significant nibble first. The entire string is preceded by the sequence \x (to distinguish it from the escape format). In some contexts, the initial backslash may need to be escaped by doubling it (see Section 4.1.2.1). For input, the hexadecimal digits can be either upper or lower case, and whitespace is permitted between digit pairs (but not within a digit pair nor in the starting \x sequence). The hex format is compatible with a wide range of external applications and protocols, and it tends to be faster to convert than the escape format, so its use is preferred.

Example:

```sql
SELECT '\xDEADBEEF';
```
8.4.2. bytea Escape Format

The “escape” format is the traditional PostgreSQL format for the bytea type. It takes the approach of representing a binary string as a sequence of ASCII characters, while converting those bytes that cannot be represented as an ASCII character into special escape sequences. If, from the point of view of the application, representing bytes as characters makes sense, then this representation can be convenient. But in practice it is usually confusing because it fuzzes up the distinction between binary strings and character strings, and also the particular escape mechanism that was chosen is somewhat unwieldy. Therefore, this format should probably be avoided for most new applications.

When entering bytea values in escape format, octets of certain values must be escaped, while all octet values can be escaped. In general, to escape an octet, convert it into its three-digit octal value and precede it by a backslash. Backslash itself (octet decimal value 92) can alternatively be represented by double backslashes. Table 8-7 shows the characters that must be escaped, and gives the alternative escape sequences where applicable.

<table>
<thead>
<tr>
<th>Decimal Octet Value</th>
<th>Description</th>
<th>Escaped Input Representation</th>
<th>Example</th>
<th>Hex Representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>zero octet</td>
<td>'\000'</td>
<td>SELECT '\000':bytea;</td>
<td>\x00</td>
</tr>
<tr>
<td>39</td>
<td>single quote</td>
<td>&quot;&quot; or '\047'</td>
<td>SELECT &quot;&quot;:bytea;</td>
<td>\x27</td>
</tr>
<tr>
<td>92</td>
<td>backslash</td>
<td>'' or '\134'</td>
<td>SELECT '':bytea;</td>
<td>\x5c</td>
</tr>
<tr>
<td>0 to 31 and 127 to 255</td>
<td>“non-printable” octets</td>
<td>'\xxx' (octal value)</td>
<td>SELECT '\001':bytea;</td>
<td>\x01</td>
</tr>
</tbody>
</table>

The requirement to escape non-printable octets varies depending on locale settings. In some instances you can get away with leaving them unescaped.

The reason that single quotes must be doubled, as shown in Table 8-7, is that this is true for any string literal in a SQL command. The generic string-literal parser consumes the outermost single quotes and reduces any pair of single quotes to one data character. What the bytea input function sees is just one single quote, which it treats as a plain data character. However, the bytea input function treats backslashes as special, and the other behaviors shown in Table 8-7 are implemented by that function.

In some contexts, backslashes must be doubled compared to what is shown above, because the generic string-literal parser will also reduce pairs of backslashes to one data character; see Section 4.1.2.1.

Bytea octets are output in hex format by default. If you change bytea_output to escape, “non-printable” octets are converted to their equivalent three-digit octal value and preceded by one backslash. Most “printable” octets are output by their standard representation in the client character set, e.g.:

SET bytea_output = 'escape';
SELECT 'abc \153\154\155 \052\251\124'::bytea;
    bytea
----------------------------------
   abc klm *\251T

The octet with decimal value 92 (backslash) is doubled in the output. Details are in Table 8-8.

Table 8-8. bytea Output Escaped Octets

<table>
<thead>
<tr>
<th>Decimal Octet Value</th>
<th>Description</th>
<th>Escaped Output Representation</th>
<th>Example</th>
<th>Output Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>92</td>
<td>backslash</td>
<td>\   \</td>
<td>SELECT '\134':bytea;</td>
<td>\</td>
</tr>
<tr>
<td>0 to 31 and 127 to 255</td>
<td>“non-printable” octets</td>
<td>\xxx (octal value)</td>
<td>SELECT '\001':bytea;</td>
<td>\001</td>
</tr>
<tr>
<td>32 to 126</td>
<td>“printable” octets</td>
<td>client character set representation</td>
<td>SELECT '\176':bytea;</td>
<td>~</td>
</tr>
</tbody>
</table>

Depending on the front end to PostgreSQL you use, you might have additional work to do in terms of escaping and unescaping bytea strings. For example, you might also have to escape line feeds and carriage returns if your interface automatically translates these.

8.5. Date/Time Types

PostgreSQL supports the full set of SQL date and time types, shown in Table 8-9. The operations available on these data types are described in Section 9.9. Dates are counted according to the Gregorian calendar, even in years before that calendar was introduced (see Section B.5 for more information).

Table 8-9. Date/Time Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Low Value</th>
<th>High Value</th>
<th>Resolution</th>
</tr>
</thead>
<tbody>
<tr>
<td>timestamp [ (p) ] [ without time zone ]</td>
<td>8 bytes</td>
<td>both date and time (no time zone)</td>
<td>4713 BC</td>
<td>294276 AD</td>
<td>1 microsecond / 14 digits</td>
</tr>
<tr>
<td>timestamp [ (p) ] with time zone</td>
<td>8 bytes</td>
<td>both date and time, with time zone</td>
<td>4713 BC</td>
<td>294276 AD</td>
<td>1 microsecond / 14 digits</td>
</tr>
<tr>
<td>date</td>
<td>4 bytes</td>
<td>date (no time of day)</td>
<td>4713 BC</td>
<td>5874897 AD</td>
<td>1 day</td>
</tr>
</tbody>
</table>
### Chapter 8. Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Low Value</th>
<th>High Value</th>
<th>Resolution</th>
</tr>
</thead>
<tbody>
<tr>
<td>time [ (p) ] [ without time zone ]</td>
<td>8 bytes</td>
<td>time of day (no date)</td>
<td>00:00:00</td>
<td>24:00:00</td>
<td>1 microsecond / 14 digits</td>
</tr>
<tr>
<td>time [ (p) ] with time zone</td>
<td>12 bytes</td>
<td>times of day only, with time zone</td>
<td>00:00:00+1459</td>
<td>24:00:00-1459</td>
<td>1 microsecond / 14 digits</td>
</tr>
<tr>
<td>interval [ fields ] [ (p) ]</td>
<td>16 bytes</td>
<td>time interval</td>
<td>-178000000 years</td>
<td>178000000 years</td>
<td>1 microsecond / 14 digits</td>
</tr>
</tbody>
</table>

**Note:** The SQL standard requires that writing just `timestamp` be equivalent to `timestamp without time zone`, and PostgreSQL honors that behavior. `timestamptz` is accepted as an abbreviation for `timestamp with time zone`; this is a PostgreSQL extension.

`time`, `timestamp`, and `interval` accept an optional precision value `p` which specifies the number of fractional digits retained in the seconds field. By default, there is no explicit bound on precision. The allowed range of `p` is from 0 to 6 for the `timestamp` and `interval` types.

**Note:** When `timestamp` values are stored as eight-byte integers (currently the default), microsecond precision is available over the full range of values. When `timestamp` values are stored as double precision floating-point numbers instead (a deprecated compile-time option), the effective limit of precision might be less than 6. `timestamp` values are stored as seconds before or after midnight 2000-01-01. When `timestamp` values are implemented using floating-point numbers, microsecond precision is achieved for dates within a few years of 2000-01-01, but the precision degrades for dates further away. Note that using floating-point datetimes allows a larger range of `timestamp` values to be represented than shown above: from 4713 BC up to 5874897 AD.

The same compile-time option also determines whether `time` and `interval` values are stored as floating-point numbers or eight-byte integers. In the floating-point case, large `interval` values degrade in precision as the size of the interval increases.

For the `time` types, the allowed range of `p` is from 0 to 6 when eight-byte integer storage is used, or from 0 to 10 when floating-point storage is used.

The `interval` type has an additional option, which is to restrict the set of stored fields by writing one of these phrases:

- `YEAR`
- `MONTH`
- `DAY`
- `HOUR`
- `MINUTE`
- `SECOND`
- `YEAR TO MONTH`
- `DAY TO HOUR`
- `DAY TO MINUTE`
- `DAY TO SECOND`
- `HOUR TO MINUTE`
- `HOUR TO SECOND`
MINUTE TO SECOND

Note that if both fields and $p$ are specified, the fields must include SECOND, since the precision applies only to the seconds.

The type time with time zone is defined by the SQL standard, but the definition exhibits properties which lead to questionable usefulness. In most cases, a combination of date, time, timestamp without time zone, and timestamp with time zone should provide a complete range of date/time functionality required by any application.

The types abstime and reltime are lower precision types which are used internally. You are discouraged from using these types in applications; these internal types might disappear in a future release.

8.5.1. Date/Time Input

Date and time input is accepted in almost any reasonable format, including ISO 8601, SQL-compatible, traditional POSTGRES, and others. For some formats, ordering of day, month, and year in date input is ambiguous and there is support for specifying the expected ordering of these fields. Set the DateStyle parameter to MDY to select month-day-year interpretation, DMY to select day-month-year interpretation, or YMD to select year-month-day interpretation.

PostgreSQL is more flexible in handling date/time input than the SQL standard requires. See Appendix B for the exact parsing rules of date/time input and for the recognized text fields including months, days of the week, and time zones.

Remember that any date or time literal input needs to be enclosed in single quotes, like text strings. Refer to Section 4.1.2.7 for more information. SQL requires the following syntax

\[ \text{type [ (p) ] 'value'} \]

where $p$ is an optional precision specification giving the number of fractional digits in the seconds field. Precision can be specified for time, timestamp, and interval types. The allowed values are mentioned above. If no precision is specified in a constant specification, it defaults to the precision of the literal value.

8.5.1.1. Dates

Table 8-10 shows some possible inputs for the date type.

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1999-01-08</td>
<td>ISO 8601; January 8 in any mode (recommended format)</td>
</tr>
<tr>
<td>January 8, 1999</td>
<td>unambiguous in any datestyle input mode</td>
</tr>
<tr>
<td>1/8/1999</td>
<td>January 8 in MDY mode; August 1 in DMY mode</td>
</tr>
<tr>
<td>1/18/1999</td>
<td>January 18 in MDY mode; rejected in other modes</td>
</tr>
<tr>
<td>01/02/03</td>
<td>January 2, 2003 in MDY mode; February 1, 2003 in DMY mode; February 3, 2001 in YMD mode</td>
</tr>
<tr>
<td>1999-Jan-08</td>
<td>January 8 in any mode</td>
</tr>
</tbody>
</table>
Chapter 8. Data Types

8.5.1.2. Times

The time-of-day types are `time [ (p) ] without time zone` and `time [ (p) ] with time zone`. `time alone` is equivalent to `time without time zone`.

Valid input for these types consists of a time of day followed by an optional time zone. (See Table 8-11 and Table 8-12.) If a time zone is specified in the input for `time without time zone`, it is silently ignored. You can also specify a date but it will be ignored, except when you use a time zone name that involves a daylight-savings rule, such as `America/New_York`. In this case specifying the date is required in order to determine whether standard or daylight-savings time applies. The appropriate time zone offset is recorded in the `time with time zone` value.

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>04:05:06.789</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05:06</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>040506</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05 AM</td>
<td>same as 04:05; AM does not affect value</td>
</tr>
<tr>
<td>04:05 PM</td>
<td>same as 16:05; input hour must be &lt;= 12</td>
</tr>
<tr>
<td>04:05:06.789-8</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05:06-08:00</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05-08:00</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>040506-08</td>
<td>ISO 8601</td>
</tr>
<tr>
<td>04:05:06 PST</td>
<td>time zone specified by abbreviation</td>
</tr>
<tr>
<td>2003-04-12 04:05:06 America/New_York</td>
<td>time zone specified by full name</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PST</td>
<td>Abbreviation (for Pacific Standard Time)</td>
</tr>
</tbody>
</table>
Chapter 8. Data Types

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>America/New_York</td>
<td>Full time zone name</td>
</tr>
<tr>
<td>PST8PDT</td>
<td>POSIX-style time zone specification</td>
</tr>
<tr>
<td>-8:00</td>
<td>ISO-8601 offset for PST</td>
</tr>
<tr>
<td>-800</td>
<td>ISO-8601 offset for PST</td>
</tr>
<tr>
<td>-8</td>
<td>ISO-8601 offset for PST</td>
</tr>
<tr>
<td>zulu</td>
<td>Military abbreviation for UTC</td>
</tr>
<tr>
<td>z</td>
<td>Short form of zulu</td>
</tr>
</tbody>
</table>

Refer to Section 8.5.3 for more information on how to specify time zones.

8.5.1.3. Time Stamps

Valid input for the time stamp types consists of the concatenation of a date and a time, followed by an optional time zone, followed by an optional AD or BC. (Alternatively, AD/BC can appear before the time zone, but this is not the preferred ordering.) Thus:

1999-01-08 04:05:06

and:

1999-01-08 04:05:06 -8:00

are valid values, which follow the ISO 8601 standard. In addition, the common format:

January 8 04:05:06 1999 PST

is supported.

The SQL standard differentiates timestamp without time zone and timestamp with time zone literals by the presence of a “+” or “-” symbol and time zone offset after the time. Hence, according to the standard,

TIMESTAMP ‘2004-10-19 10:23:54’

is a timestamp without time zone, while

TIMESTAMP ‘2004-10-19 10:23:54+02’

is a timestamp with time zone. PostgreSQL never examines the content of a literal string before determining its type, and therefore will treat both of the above as timestamp without time zone. To ensure that a literal is treated as timestamp with time zone, give it the correct explicit type:

TIMESTAMP WITH TIME ZONE ‘2004-10-19 10:23:54+02’

In a literal that has been determined to be timestamp without time zone, PostgreSQL will silently ignore any time zone indication. That is, the resulting value is derived from the date/time fields in the input value, and is not adjusted for time zone.

For timestamp with time zone, the internally stored value is always in UTC (Universal Coordinated Time, traditionally known as Greenwich Mean Time, GMT). An input value that has an explicit time zone specified is converted to UTC using the appropriate offset for that time zone. If no time zone is stated in the input string, then it is assumed to be in the time zone indicated by the system’s TimeZone parameter, and is converted to UTC using the offset for the timezone zone.
When a timestamp with time zone value is output, it is always converted from UTC to the current timezone zone, and displayed as local time in that zone. To see the time in another time zone, either change timezone or use the AT TIME ZONE construct (see Section 9.9.3).

Conversions between timestamp without time zone and timestamp with time zone normally assume that the timestamp without time zone value should be taken or given as timezone local time. A different time zone can be specified for the conversion using AT TIME ZONE.

### 8.5.1.4. Special Values

PostgreSQL supports several special date/time input values for convenience, as shown in Table 8-13. The values infinity and -infinity are specially represented inside the system and will be displayed unchanged; but the others are simply notational shorthands that will be converted to ordinary date/time values when read. (In particular, now and related strings are converted to a specific time value as soon as they are read.) All of these values need to be enclosed in single quotes when used as constants in SQL commands.

<table>
<thead>
<tr>
<th>Input String</th>
<th>Valid Types</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>epoch</td>
<td>date,timestamp</td>
<td>1970-01-01 00:00:00+00 (Unix system time zero)</td>
</tr>
<tr>
<td>infinity</td>
<td>date,timestamp</td>
<td>later than all other time stamps</td>
</tr>
<tr>
<td>-infinity</td>
<td>date,timestamp</td>
<td>earlier than all other time stamps</td>
</tr>
<tr>
<td>now</td>
<td>date,time,timestamp</td>
<td>current transaction’s start time</td>
</tr>
<tr>
<td>today</td>
<td>date,timestamp</td>
<td>midnight (00:00) today</td>
</tr>
<tr>
<td>tomorrow</td>
<td>date,timestamp</td>
<td>midnight (00:00) tomorrow</td>
</tr>
<tr>
<td>yesterday</td>
<td>date,timestamp</td>
<td>midnight (00:00) yesterday</td>
</tr>
<tr>
<td>allballs</td>
<td>time</td>
<td>00:00:00.00 UTC</td>
</tr>
</tbody>
</table>

The following SQL-compatible functions can also be used to obtain the current time value for the corresponding data type: CURRENT_DATE, CURRENT_TIME, CURRENT_TIMESTAMP, LOCALTIME, LOCALTIMESTAMP. The latter four accept an optional subsecond precision specification. (See Section 9.9.4.) Note that these are SQL functions and are not recognized in data input strings.

### 8.5.2. Date/Time Output

The output format of the date/time types can be set to one of the four styles ISO 8601, SQL (Ingres), traditional POSTGRES (Unix date format), or German. The default is the ISO format. (The SQL standard requires the use of the ISO 8601 format. The name of the “SQL” output format is a historical accident.) Table 8-14 shows examples of each output style. The output of the date and time types is generally only the date or time part in accordance with the given examples. However, the POSTGRES style outputs date-only values in ISO format.

<table>
<thead>
<tr>
<th>Input String</th>
<th>Valid Types</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>date</td>
<td>midnight (00:00) today</td>
</tr>
<tr>
<td></td>
<td>time</td>
<td>00:00:00.00 UTC</td>
</tr>
</tbody>
</table>

**Chapter 8. Data Types**
Chapter 8. Data Types

<table>
<thead>
<tr>
<th>Style Specification</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>ISO</td>
<td>ISO 8601, SQL standard</td>
<td>1997-12-17 07:37:16-08</td>
</tr>
<tr>
<td>SQL</td>
<td>traditional style</td>
<td>12/17/1997 07:37:16.00 PST</td>
</tr>
<tr>
<td>Postgres</td>
<td>original style</td>
<td>Wed Dec 17 07:37:16 1997 PST</td>
</tr>
<tr>
<td>German</td>
<td>regional style</td>
<td>17.12.1997 07:37:16.00 PST</td>
</tr>
</tbody>
</table>

Note: ISO 8601 specifies the use of uppercase letter T to separate the date and time. PostgreSQL accepts that format on input, but on output it uses a space rather than T, as shown above. This is for readability and for consistency with RFC 3339 as well as some other database systems.

In the SQL and POSTGRES styles, day appears before month if DMY field ordering has been specified, otherwise month appears before day. (See Section 8.5.1 for how this setting also affects interpretation of input values.) Table 8-15 shows examples.

Table 8-15. Date Order Conventions

<table>
<thead>
<tr>
<th>datestyle Setting</th>
<th>Input Ordering</th>
<th>Example Output</th>
</tr>
</thead>
<tbody>
<tr>
<td>SQL, DMY</td>
<td>day/month/year</td>
<td>17/12/1997 15:37:16.00 CET</td>
</tr>
<tr>
<td>SQL, MDY</td>
<td>month/day/year</td>
<td>12/17/1997 07:37:16.00 PST</td>
</tr>
<tr>
<td>Postgres, DMY</td>
<td>day/month/year</td>
<td>Wed 17 Dec 07:37:16 1997 PST</td>
</tr>
</tbody>
</table>

The date/time style can be selected by the user using the SET datestyle command, the DateStyle parameter in the postgresql.conf configuration file, or the PGDATESTYLE environment variable on the server or client.

The formatting function to_char (see Section 9.8) is also available as a more flexible way to format date/time output.

8.5.3. Time Zones

Time zones, and time-zone conventions, are influenced by political decisions, not just earth geometry. Time zones around the world became somewhat standardized during the 1900s, but continue to be prone to arbitrary changes, particularly with respect to daylight-savings rules. PostgreSQL uses the widely-used IANA (Olson) time zone database for information about historical time zone rules. For times in the future, the assumption is that the latest known rules for a given time zone will continue to be observed indefinitely far into the future.

PostgreSQL endeavors to be compatible with the SQL standard definitions for typical usage. However, the SQL standard has an odd mix of date and time types and capabilities. Two obvious problems are:

- Although the date type cannot have an associated time zone, the time type can. Time zones in the real world have little meaning unless associated with a date as well as a time, since the offset can vary through the year with daylight-saving time boundaries.
The default time zone is specified as a constant numeric offset from UTC. It is therefore impossible to adapt to daylight-saving time when doing date/time arithmetic across DST boundaries.

To address these difficulties, we recommend using date/time types that contain both date and time when using time zones. We do not recommend using the type `time with time zone` (though it is supported by PostgreSQL for legacy applications and for compliance with the SQL standard). PostgreSQL assumes your local time zone for any type containing only date or time.

All timezone-aware dates and times are stored internally in UTC. They are converted to local time in the zone specified by the TimeZone configuration parameter before being displayed to the client.

PostgreSQL allows you to specify time zones in three different forms:

- A full time zone name, for example `America/New_York`. The recognized time zone names are listed in the `pg_timezone_names` view (see Section 50.80). PostgreSQL uses the widely-used IANA time zone data for this purpose, so the same time zone names are also recognized by other software.

- A time zone abbreviation, for example `PST`. Such a specification merely defines a particular offset from UTC, in contrast to full time zone names which can imply a set of daylight savings transition-date rules as well. The recognized abbreviations are listed in the `pg_timezone_abbrevs` view (see Section 50.79). You cannot set the configuration parameters TimeZone or log_timezone to a time zone abbreviation, but you can use abbreviations in date/time input values and with the `AT TIME ZONE` operator.

- In addition to the timezone names and abbreviations, PostgreSQL will accept POSIX-style time zone specifications of the form `STDoffset` or `STDoffsetDST`, where `STD` is a zone abbreviation, `offset` is a numeric offset in hours west from UTC, and `DST` is an optional daylight-savings zone abbreviation, assumed to stand for one hour ahead of the given offset. For example, if `ESTSEDT` were not already a recognized zone name, it would be accepted and would be functionally equivalent to United States East Coast time. In this syntax, a zone abbreviation can be a string of letters, or an arbitrary string surrounded by angle brackets (<>). When a daylight-savings zone abbreviation is present, it is assumed to be used according to the same daylight-savings transition rules used in the IANA time zone database’s `posixrules` entry. In a standard PostgreSQL installation, `posixrules` is the same as `US/Eastern`, so that POSIX-style time zone specifications follow USA daylight-savings rules. If needed, you can adjust this behavior by replacing the `posixrules` file.

In short, this is the difference between abbreviations and full names: abbreviations represent a specific offset from UTC, whereas many of the full names imply a local daylight-savings time rule, and so have two possible UTC offsets. As an example, `2014-06-04 12:00 America/New_York` represents noon local time in New York, which for this particular date was Eastern Daylight Time (UTC-4). So `2014-06-04 12:00 EDT` specifies that same time instant. But `2014-06-04 12:00 EST` specifies noon Eastern Standard Time (UTC-5), regardless of whether daylight savings was nominally in effect on that date.

To complicate matters, some jurisdictions have used the same timezone abbreviation to mean different UTC offsets at different times; for example, in Moscow `MSK` has meant UTC+3 in some years and UTC+4 in others. PostgreSQL interprets such abbreviations according to whatever they meant (or had most recently meant) on the specified date; but, as with the `EST` example above, this is not necessarily the same as local civil time on that date.

One should be wary that the POSIX-style time zone feature can lead to silently accepting bogus input, since there is no check on the reasonableness of the zone abbreviations. For example, `SET TIMEZONE TO FOOBAR0` will work, leaving the system effectively using a rather peculiar abbreviation for UTC.
Another issue to keep in mind is that in POSIX time zone names, positive offsets are used for locations west of Greenwich. Everywhere else, PostgreSQL follows the ISO-8601 convention that positive timezone offsets are east of Greenwich.

In all cases, timezone names and abbreviations are recognized case-insensitively. (This is a change from PostgreSQL versions prior to 8.2, which were case-sensitive in some contexts but not others.) Neither timezone names nor abbreviations are hard-wired into the server; they are obtained from configuration files stored under .../share/timezone/ and .../share/timezonesets/ of the installation directory (see Section B.4).

The TimeZone configuration parameter can be set in the file postgresql.conf, or in any of the other standard ways described in Chapter 19. There are also some special ways to set it:

- The SQL command SET TIME ZONE sets the time zone for the session. This is an alternative spelling of SET TIMEZONE TO with a more SQL-spec-compatible syntax.
- The PGTZ environment variable is used by libpq clients to send a SET TIME ZONE command to the server upon connection.

### 8.5.4. Interval Input

Interval values can be written using the following verbose syntax:

\[
[@] \text{quantity unit } [\text{quantity unit}... ] [\text{direction}]
\]

where quantity is a number (possibly signed); unit is microsecond, millisecond, second, minute, hour, day, week, month, year, decade, century, millennium, or abbreviations or plurals of these units; direction can be ago or empty. The at sign (@) is optional noise. The amounts of the different units are implicitly added with appropriate sign accounting. ago negates all the fields. This syntax is also used for interval output, if IntervalStyle is set to postgres_verbose.

Quantities of days, hours, minutes, and seconds can be specified without explicit unit markings. For example, ‘1 12:59:10’ is read the same as ‘1 day 12 hours 59 min 10 sec’. Also, a combination of years and months can be specified with a dash; for example ‘200-10’ is read the same as ‘200 years 10 months’. (These shorter forms are in fact the only ones allowed by the SQL standard, and are used for output when IntervalStyle is set to sql_standard.)

Interval values can also be written as ISO 8601 time intervals, using either the “format with designators” of the standard’s section 4.4.3.2 or the “alternative format” of section 4.4.3.3. The format with designators looks like this:

\[
P \text{quantity unit } [ \text{quantity unit}... ] [ \text{T } [\text{quantity unit }... ]]
\]

The string must start with a P, and may include a T that introduces the time-of-day units. The available unit abbreviations are given in Table 8-16. Units may be omitted, and may be specified in any order, but units smaller than a day must appear after T. In particular, the meaning of M depends on whether it is before or after T.

### Table 8-16. ISO 8601 Interval Unit Abbreviations

<table>
<thead>
<tr>
<th>Abbreviation</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>Y</td>
<td>Years</td>
</tr>
</tbody>
</table>
### Abbreviation

<table>
<thead>
<tr>
<th>Abbreviation</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>M</td>
<td>Months (in the date part)</td>
</tr>
<tr>
<td>W</td>
<td>Weeks</td>
</tr>
<tr>
<td>D</td>
<td>Days</td>
</tr>
<tr>
<td>H</td>
<td>Hours</td>
</tr>
<tr>
<td>M</td>
<td>Minutes (in the time part)</td>
</tr>
<tr>
<td>S</td>
<td>Seconds</td>
</tr>
</tbody>
</table>

In the alternative format:

\[
P \ [ \text{years-months-days} \ ] \ [ T \text{hours:minutes:seconds} ]
\]

the string must begin with \( P \), and a \( T \) separates the date and time parts of the interval. The values are given as numbers similar to ISO 8601 dates.

When writing an interval constant with a \texttt{fields} specification, or when assigning a string to an interval column that was defined with a \texttt{fields} specification, the interpretation of unmarked quantities depends on the \texttt{fields}. For example \texttt{INTERVAL ‘1’ YEAR} is read as 1 year, whereas \texttt{INTERVAL ‘1’} means 1 second. Also, field values “to the right” of the least significant field allowed by the \texttt{fields} specification are silently discarded. For example, writing \texttt{INTERVAL ‘1 day 2:03:04’ HOUR TO MINUTE} results in dropping the seconds field, but not the day field.

According to the SQL standard all fields of an interval value must have the same sign, so a leading negative sign applies to all fields; for example the negative sign in the interval literal \( ‘-1 2:03:04’ \) applies to both the days and hour/minute/second parts. PostgreSQL allows the fields to have different signs, and traditionally treats each field in the textual representation as independently signed, so that the hour/minute/second part is considered positive in this example. If \texttt{IntervalStyle} is set to \texttt{sql_standard} then a leading sign is considered to apply to all fields (but only if no additional signs appear). Otherwise the traditional PostgreSQL interpretation is used. To avoid ambiguity, it’s recommended to attach an explicit sign to each field if any field is negative.

In the verbose input format, and in some fields of the more compact input formats, field values can have fractional parts; for example \( ‘1.5 week’ \) or \( ‘01:02:03.45’ \). Such input is converted to the appropriate number of months, days, and seconds for storage. When this would result in a fractional number of months or days, the fraction is added to the lower-order fields using the conversion factors 1 month = 30 days and 1 day = 24 hours. For example, \( ‘1.5 month’ \) becomes 1 month and 15 days. Only seconds will ever be shown as fractional on output.

Table 8-17 shows some examples of valid interval input.

### Table 8-17. Interval Input

<table>
<thead>
<tr>
<th>Example</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1-2</td>
<td>SQL standard format: 1 year 2 months</td>
</tr>
<tr>
<td>3 4:05:06</td>
<td>SQL standard format: 3 days 4 hours 5 minutes 6 seconds</td>
</tr>
<tr>
<td>1 year 2 months 3 days 4 hours 5 minutes 6 seconds</td>
<td>Traditional Postgres format: 1 year 2 months 3 days 4 hours 5 minutes 6 seconds</td>
</tr>
<tr>
<td>P1Y2M3DT4H5M6S</td>
<td>ISO 8601 “format with designators”: same meaning as above</td>
</tr>
<tr>
<td>P0001-02-03T04:05:06</td>
<td>ISO 8601 “alternative format”: same meaning as above</td>
</tr>
</tbody>
</table>
Internally, interval values are stored as months, days, and seconds. This is done because the number of days in a month varies, and a day can have 23 or 25 hours if a daylight savings time adjustment is involved. The months and days fields are integers while the seconds field can store fractions. Because intervals are usually created from constant strings or timestamp subtraction, this storage method works well in most cases, but can cause unexpected results:

```
SELECT EXTRACT(hours from '80 minutes '::interval);
          date_part
-------------
             1
```

```
SELECT EXTRACT(days from '80 hours '::interval);
          date_part
-------------
             0
```

Functions `justify_days` and `justify_hours` are available for adjusting days and hours that overflow their normal ranges.

### 8.5.5. Interval Output

The output format of the interval type can be set to one of the four styles `sql_standard`, `postgres`, `postgres_verbose`, or `iso_8601`, using the command `SET intervalstyle`. The default is the `postgres` format. Table 8-18 shows examples of each output style.

The `sql_standard` style produces output that conforms to the SQL standard’s specification for interval literal strings, if the interval value meets the standard’s restrictions (either year-month only or day-time only, with no mixing of positive and negative components). Otherwise the output looks like a standard year-month literal string followed by a day-time literal string, with explicit signs added to disambiguate mixed-sign intervals.

The output of the `postgres` style matches the output of PostgreSQL releases prior to 8.4 when the `DateStyle` parameter was set to `ISO`.

The output of the `postgres_verbose` style matches the output of PostgreSQL releases prior to 8.4 when the `DateStyle` parameter was set to non-ISO output.

The output of the `iso_8601` style matches the “format with designators” described in section 4.4.3.2 of the ISO 8601 standard.

<table>
<thead>
<tr>
<th>Style Specification</th>
<th>Year-Month Interval</th>
<th>Day-Time Interval</th>
<th>Mixed Interval</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>sql_standard</code></td>
<td>1-2</td>
<td>3 4:05:06</td>
<td>-1-2 +3 -4:05:06</td>
</tr>
<tr>
<td><code>postgres</code></td>
<td>1 year 2 mons</td>
<td>3 days 04:05:06</td>
<td>-1 year -2 mons +3 days -04:05:06</td>
</tr>
<tr>
<td><code>postgres_verbose</code></td>
<td>@ 1 year 2 mons</td>
<td>@ 3 days 4 hours 5 mins 6 secs</td>
<td>@ 1 year 2 mons -3 days 4 hours 5 mins 6 secs ago</td>
</tr>
<tr>
<td><code>iso_8601</code></td>
<td>P1Y2M</td>
<td>P3DT4H5M6S</td>
<td>P-1Y-2M3DT-4H-5M-6S</td>
</tr>
</tbody>
</table>
8.6. Boolean Type

PostgreSQL provides the standard SQL type `boolean`; see Table 8-19. The `boolean` type can have several states: “true”, “false”, and a third state, “unknown”, which is represented by the SQL null value.

Table 8-19. Boolean Data Type

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>boolean</td>
<td>1 byte</td>
<td>state of true or false</td>
</tr>
</tbody>
</table>

Boolean constants can be represented in SQL queries by the SQL key words `TRUE`, `FALSE`, and `NULL`. The datatype input function for type `boolean` accepts these string representations for the “true” state:

```
true
yes
on
1
```

and these representations for the “false” state:

```
false
no
off
0
```

Unique prefixes of these strings are also accepted, for example `t` or `n`. Leading or trailing whitespace is ignored, and case does not matter.

The datatype output function for type `boolean` always emits either `t` or `f`, as shown in Example 8-2.

Example 8-2. Using the boolean Type

```sql
CREATE TABLE test1 (a boolean, b text);
INSERT INTO test1 VALUES (TRUE, 'sic est');
INSERT INTO test1 VALUES (FALSE, 'non est');
SELECT * FROM test1;
  a | b
---+---------
  t | sic est
  f | non est

SELECT * FROM test1 WHERE a;
  a | b
---+---------
  t | sic est
```

The key words `TRUE` and `FALSE` are the preferred (SQL-compliant) method for writing Boolean constants in SQL queries. But you can also use the string representations by following the generic string-literal constant syntax described in Section 4.1.2.7, for example `'yes'::boolean`.

Note that the parser automatically understands that `TRUE` and `FALSE` are of type `boolean`, but this is not so for `NULL` because that can have any type. So in some contexts you might have to cast `NULL` to `boolean` explicitly, for example `NULL::boolean`. Conversely, the cast can be omitted from a
string-literal Boolean value in contexts where the parser can deduce that the literal must be of type `boolean`.

### 8.7. Enumerated Types

Enumerated (enum) types are data types that comprise a static, ordered set of values. They are equivalent to the `enum` types supported in a number of programming languages. An example of an enum type might be the days of the week, or a set of status values for a piece of data.

#### 8.7.1. Declaration of Enumerated Types

Enum types are created using the `CREATE TYPE` command, for example:

```sql
CREATE TYPE mood AS ENUM ('sad', 'ok', 'happy');
```

Once created, the enum type can be used in table and function definitions much like any other type:

```sql
CREATE TABLE person (  
  name text,  
  current_mood mood
);
```

```sql
INSERT INTO person VALUES ('Moe', 'happy');
SELECT * FROM person WHERE current_mood = 'happy';
```

<table>
<thead>
<tr>
<th>name</th>
<th>current_mood</th>
</tr>
</thead>
<tbody>
<tr>
<td>Moe</td>
<td>happy</td>
</tr>
</tbody>
</table>

(1 row)

#### 8.7.2. Ordering

The ordering of the values in an enum type is the order in which the values were listed when the type was created. All standard comparison operators and related aggregate functions are supported for enums. For example:

```sql
INSERT INTO person VALUES ('Larry', 'sad');
INSERT INTO person VALUES ('Curly', 'ok');
SELECT * FROM person WHERE current_mood > 'sad';
```

<table>
<thead>
<tr>
<th>name</th>
<th>current_mood</th>
</tr>
</thead>
<tbody>
<tr>
<td>Moe</td>
<td>happy</td>
</tr>
<tr>
<td>Curly</td>
<td>ok</td>
</tr>
</tbody>
</table>

(2 rows)

```sql
SELECT * FROM person WHERE current_mood > 'sad' ORDER BY current_mood;
```

<table>
<thead>
<tr>
<th>name</th>
<th>current_mood</th>
</tr>
</thead>
<tbody>
<tr>
<td>Curly</td>
<td>ok</td>
</tr>
<tr>
<td>Moe</td>
<td>happy</td>
</tr>
</tbody>
</table>

(2 rows)
Chapter 8. Data Types

SELECT name
FROM person
WHERE current_mood = (SELECT MIN(current_mood) FROM person);

<table>
<thead>
<tr>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>Larry</td>
</tr>
</tbody>
</table>

(1 row)

8.7.3. Type Safety

Each enumerated data type is separate and cannot be compared with other enumerated types. See this example:

CREATE TYPE happiness AS ENUM ('happy', 'very happy', 'ecstatic');
CREATE TABLE holidays (num_weeks integer, happiness happiness);
INSERT INTO holidays(num_weeks,happiness) VALUES (4, 'happy');
INSERT INTO holidays(num_weeks,happiness) VALUES (6, 'very happy');
INSERT INTO holidays(num_weeks,happiness) VALUES (8, 'ecstatic');
INSERT INTO holidays(num_weeks,happiness) VALUES (2, 'sad');
ERROR: invalid input value for enum happiness: "sad"
SELECT person.name, holidays.num_weeks FROM person, holidays
WHERE person.current_mood = holidays.happiness;
ERROR: operator does not exist: mood = happiness

If you really need to do something like that, you can either write a custom operator or add explicit casts to your query:

SELECT person.name, holidays.num_weeks FROM person, holidays
WHERE person.current_mood::text = holidays.happiness::text;

<table>
<thead>
<tr>
<th>name</th>
<th>num_weeks</th>
</tr>
</thead>
<tbody>
<tr>
<td>Moe</td>
<td>4</td>
</tr>
</tbody>
</table>

(1 row)

8.7.4. Implementation Details

Enum labels are case sensitive, so 'happy' is not the same as 'HAPPY'. White space in the labels is significant too.

Although enum types are primarily intended for static sets of values, there is support for adding new values to an existing enum type, and for renaming values (see ALTER TYPE). Existing values cannot be removed from an enum type, nor can the sort ordering of such values be changed, short of dropping and re-creating the enum type.
An enum value occupies four bytes on disk. The length of an enum value’s textual label is limited by the `NAMEDATALEN` setting compiled into PostgreSQL; in standard builds this means at most 63 bytes. The translations from internal enum values to textual labels are kept in the system catalog `pg_enum`. Querying this catalog directly can be useful.

### 8.8. Geometric Types

Geometric data types represent two-dimensional spatial objects. Table 8-20 shows the geometric types available in PostgreSQL.

#### Table 8-20. Geometric Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
<th>Representation</th>
</tr>
</thead>
<tbody>
<tr>
<td>point</td>
<td>16 bytes</td>
<td>Point on a plane</td>
<td>(x,y)</td>
</tr>
<tr>
<td>line</td>
<td>32 bytes</td>
<td>Infinite line</td>
<td>{A,B,C}</td>
</tr>
<tr>
<td>lseg</td>
<td>32 bytes</td>
<td>Finite line segment</td>
<td>((x1,y1),(x2,y2))</td>
</tr>
<tr>
<td>box</td>
<td>32 bytes</td>
<td>Rectangular box</td>
<td>((x1,y1),(x2,y2))</td>
</tr>
<tr>
<td>path</td>
<td>16+16n bytes</td>
<td>Closed path (similar to polygon)</td>
<td>((x1,y1),...)</td>
</tr>
<tr>
<td>path</td>
<td>16+16n bytes</td>
<td>Open path</td>
<td>[(x1,y1),...]</td>
</tr>
<tr>
<td>polygon</td>
<td>40+16n bytes</td>
<td>Polygon (similar to closed path)</td>
<td>((x1,y1),...)</td>
</tr>
<tr>
<td>circle</td>
<td>24 bytes</td>
<td>Circle</td>
<td>(&lt;(x,y),r&gt; (center point and radius)</td>
</tr>
</tbody>
</table>

A rich set of functions and operators is available to perform various geometric operations such as scaling, translation, rotation, and determining intersections. They are explained in Section 9.11.

#### 8.8.1. Points

Points are the fundamental two-dimensional building block for geometric types. Values of type `point` are specified using either of the following syntaxes:

```
(x, y)
```

where \(x\) and \(y\) are the respective coordinates, as floating-point numbers.

Points are output using the first syntax.

#### 8.8.2. Lines

Lines are represented by the linear equation \(Ax + By + C = 0\), where \(A\) and \(B\) are not both zero. Values of type `line` are input and output in the following form:

```
{ A, B, C }
```

Alternatively, any of the following forms can be used for input:
where \((x_1, y_1)\) and \((x_2, y_2)\) are two different points on the line.

### 8.8.3. Line Segments

Line segments are represented by pairs of points that are the endpoints of the segment. Values of type `lseg` are specified using any of the following syntaxes:

\[
[ (x_1, y_1), (x_2, y_2) ]
\]

\[
( (x_1, y_1), (x_2, y_2) )
\]

\[
( x_1, y_1, x_2, y_2 )
\]

where \((x_1, y_1)\) and \((x_2, y_2)\) are the end points of the line segment.

Line segments are output using the first syntax.

### 8.8.4. Boxes

Boxes are represented by pairs of points that are opposite corners of the box. Values of type `box` are specified using any of the following syntaxes:

\[
( (x_1, y_1), (x_2, y_2) )
\]

\[
( x_1, y_1, x_2, y_2 )
\]

where \((x_1, y_1)\) and \((x_2, y_2)\) are any two opposite corners of the box.

Boxes are output using the second syntax.

Any two opposite corners can be supplied on input, but the values will be reordered as needed to store the upper right and lower left corners, in that order.

### 8.8.5. Paths

Paths are represented by lists of connected points. Paths can be open, where the first and last points in the list are considered not connected, or closed, where the first and last points are considered connected.

Values of type `path` are specified using any of the following syntaxes:

\[
[ (x_1, y_1), \ldots, (x_n, y_n) ]
\]

\[
( (x_1, y_1), \ldots, (x_n, y_n) )
\]

\[
( x_1, y_1, \ldots, x_n, y_n )
\]

\[
( x_1, y_1, x_2, y_2 )
\]

\[
( x_1, y_1, \ldots, x_n, y_n )
\]

where the points are the end points of the line segments comprising the path. Square brackets `[[]]` indicate an open path, while parentheses `(() )` indicate a closed path. When the outermost parentheses are omitted, as in the third through fifth syntaxes, a closed path is assumed.
Chapter 8. Data Types

Paths are output using the first or second syntax, as appropriate.

8.8.6. Polygons

Polygons are represented by lists of points (the vertexes of the polygon). Polygons are very similar to closed paths, but are stored differently and have their own set of support routines.

Values of type polygon are specified using any of the following syntaxes:

\[
( ( \ x_1 \ , \ y_1 \ ) \ , \ ... \ , \ ( \ x_n \ , \ y_n \ ) )
\]
\[
( \ x_1 \ , \ y_1 \ ) \ , \ ... \ , \ ( \ x_n \ , \ y_n \ )
\]
\[
( \ x_1 \ , \ y_1 \ , \ ... \ , \ x_n \ , \ y_n \ )
\]
\[
\ x_1 \ , \ y_1 \ , \ ... \ , \ x_n \ , \ y_n
\]

where the points are the end points of the line segments comprising the boundary of the polygon.

Polygons are output using the first syntax.

8.8.7. Circles

Circles are represented by a center point and radius. Values of type circle are specified using any of the following syntaxes:

\[
< ( \ x \ , \ y \ ) \ , \ r >
\]
\[
( ( \ x \ , \ y \ ) \ , \ r )
\]
\[
( \ x \ , \ y \ ) \ , \ r
\]
\[
\ x \ , \ y \ , \ r
\]

where \((x,y)\) is the center point and \(r\) is the radius of the circle.

Circles are output using the first syntax.

8.9. Network Address Types

PostgreSQL offers data types to store IPv4, IPv6, and MAC addresses, as shown in Table 8-21. It is better to use these types instead of plain text types to store network addresses, because these types offer input error checking and specialized operators and functions (see Section 9.12).

Table 8-21. Network Address Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Storage Size</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>cidr</td>
<td>7 or 19 bytes</td>
<td>IPv4 and IPv6 networks</td>
</tr>
<tr>
<td>inet</td>
<td>7 or 19 bytes</td>
<td>IPv4 and IPv6 hosts and networks</td>
</tr>
<tr>
<td>macaddr</td>
<td>6 bytes</td>
<td>MAC addresses</td>
</tr>
</tbody>
</table>

When sorting inet or cidr data types, IPv4 addresses will always sort before IPv6 addresses, including IPv4 addresses encapsulated or mapped to IPv6 addresses, such as ::10.2.3.4 or ::ffff:10.4.3.2.
8.9.1. inet

The inet type holds an IPv4 or IPv6 host address, and optionally its subnet, all in one field. The subnet is represented by the number of network address bits present in the host address (the “netmask”). If the netmask is 32 and the address is IPv4, then the value does not indicate a subnet, only a single host. In IPv6, the address length is 128 bits, so 128 bits specify a unique host address. Note that if you want to accept only networks, you should use the cidr type rather than inet.

The input format for this type is address/y where address is an IPv4 or IPv6 address and y is the number of bits in the netmask. If the /y portion is missing, the netmask is 32 for IPv4 and 128 for IPv6, so the value represents just a single host. On display, the /y portion is suppressed if the netmask specifies a single host.

8.9.2. cidr

The cidr type holds an IPv4 or IPv6 network specification. Input and output formats follow Classless Internet Domain Routing conventions. The format for specifying networks is address/y where address is the network represented as an IPv4 or IPv6 address, and y is the number of bits in the netmask. If y is omitted, it is calculated using assumptions from the older classful network numbering system, except it will be at least large enough to include all of the octets written in the input. It is an error to specify a network address that has bits set to the right of the specified netmask.

Table 8-22 shows some examples.

Table 8-22. cidr Type Input Examples

<table>
<thead>
<tr>
<th>cidr Input</th>
<th>cidr Output</th>
<th>abbrev(cidr)</th>
</tr>
</thead>
<tbody>
<tr>
<td>192.168.100.128/25</td>
<td>192.168.100.128/25</td>
<td>192.168.100.128/25</td>
</tr>
<tr>
<td>192.168/24</td>
<td>192.168.0.0/24</td>
<td>192.168.0/24</td>
</tr>
<tr>
<td>192.168/25</td>
<td>192.168.0.0/25</td>
<td>192.168.0/25</td>
</tr>
<tr>
<td>192.168.1</td>
<td>192.168.1.0/24</td>
<td>192.168.1/24</td>
</tr>
<tr>
<td>192.168</td>
<td>192.168.0.0/24</td>
<td>192.168.0/24</td>
</tr>
<tr>
<td>128.1</td>
<td>128.1.0.0/16</td>
<td>128.1/16</td>
</tr>
<tr>
<td>128</td>
<td>128.0.0.0/16</td>
<td>128.0/16</td>
</tr>
<tr>
<td>128.1.2</td>
<td>128.1.2.0/24</td>
<td>128.1.2/24</td>
</tr>
<tr>
<td>10.1.2</td>
<td>10.1.2.0/24</td>
<td>10.1.2/24</td>
</tr>
<tr>
<td>10.1</td>
<td>10.1.0.0/16</td>
<td>10.1/16</td>
</tr>
<tr>
<td>10</td>
<td>10.0.0.0/8</td>
<td>10/8</td>
</tr>
<tr>
<td>10.1.2.3/32</td>
<td>10.1.2.3/32</td>
<td>10.1.2.3/32</td>
</tr>
<tr>
<td>2001:4f8:3:ba::/64</td>
<td>2001:4f8:3:ba::/64</td>
<td>2001:4f8:3:ba::/64</td>
</tr>
<tr>
<td>::ffff:1.2.3.0/120</td>
<td>::ffff:1.2.3.0/120</td>
<td>::ffff:1.2.3/120</td>
</tr>
<tr>
<td>::ffff:1.2.3.0/128</td>
<td>::ffff:1.2.3.0/128</td>
<td>::ffff:1.2.3/128</td>
</tr>
</tbody>
</table>

8.9.3. inet vs. cidr

The essential difference between inet and cidr data types is that inet accepts values with nonzero
bits to the right of the netmask, whereas `cidr` does not.

**Tip:** If you do not like the output format for `inet` or `cidr` values, try the functions `host`, `text`, and `abbrev`.

### 8.9.4. macaddr

The `macaddr` type stores MAC addresses, known for example from Ethernet card hardware addresses (although MAC addresses are used for other purposes as well). Input is accepted in the following formats:

- `'08:00:2b:01:02:03'`
- `'08-00-2b-01-02-03'`
- `'08002b:010203'`
- `'0800b-010203'`
- `'0800.2b01.0203'`
- `'0800-2b01-0203'`
- `'08002b010203'`

These examples would all specify the same address. Upper and lower case is accepted for the digits `a` through `f`. Output is always in the first of the forms shown.

IEEE Std 802-2001 specifies the second shown form (with hyphens) as the canonical form for MAC addresses, and specifies the first form (with colons) as the bit-reversed notation, so that `08-00-2b-01-02-03 = 01:00:4D:08:04:0C`. This convention is widely ignored nowadays, and it is relevant only for obsolete network protocols (such as Token Ring). PostgreSQL makes no provisions for bit reversal, and all accepted formats use the canonical LSB order.

The remaining five input formats are not part of any standard.

### 8.10. Bit String Types

Bit strings are strings of 1’s and 0’s. They can be used to store or visualize bit masks. There are two SQL bit types: `bit(n)` and `bit varying(n)`, where `n` is a positive integer.

`bit` type data must match the length `n` exactly; it is an error to attempt to store shorter or longer bit strings. `bit varying` data is of variable length up to the maximum length `n`; longer strings will be rejected. Writing `bit` without a length is equivalent to `bit(1)`, while `bit varying` without a length specification means unlimited length.

**Note:** If one explicitly casts a bit-string value to `bit(n)`, it will be truncated or zero-padded on the right to be exactly `n` bits, without raising an error. Similarly, if one explicitly casts a bit-string value to `bit varying(n)`, it will be truncated on the right if it is more than `n` bits.

Refer to Section 4.1.2.5 for information about the syntax of bit string constants. Bit-logical operators and string manipulation functions are available; see Section 9.6.
Example 8-3. Using the Bit String Types

CREATE TABLE test (a BIT(3), b BIT VARYING(5));
INSERT INTO test VALUES (B'101', B'00');
INSERT INTO test VALUES (B'10', B'101');
ERROR: bit string length 2 does not match type bit(3)
INSERT INTO test VALUES (B'10''::bit(3), B'101');
SELECT * FROM test;
da | b
---+-----
101 | 00
100 | 101

A bit string value requires 1 byte for each group of 8 bits, plus 5 or 8 bytes overhead depending on the length of the string (but long values may be compressed or moved out-of-line, as explained in Section 8.3 for character strings).

8.11. Text Search Types

PostgreSQL provides two data types that are designed to support full text search, which is the activity of searching through a collection of natural-language documents to locate those that best match a query. The tsvector type represents a document in a form optimized for text search; the tsquery type similarly represents a text query. Chapter 12 provides a detailed explanation of this facility, and Section 9.13 summarizes the related functions and operators.

8.11.1. tsvector

A tsvector value is a sorted list of distinct lexemes, which are words that have been normalized to merge different variants of the same word (see Chapter 12 for details). Sorting and duplicate-elimination are done automatically during input, as shown in this example:

```
SELECT 'a fat cat sat on a mat and ate a fat rat'::tsvector;
```
```
tsvectoR
------------------------------
'a' 'and' 'ate' 'cat' 'fat' 'mat' 'on' 'rat' 'sat'
```

To represent lexemes containing whitespace or punctuation, surround them with quotes:

```
SELECT $$the lexeme ' ' contains spaces$$::tsvector;
```
```
tsvectoR
-------------------------------------------
' ' 'contains' 'lexeme' 'spaces' 'the'
```

(We use dollar-quoted string literals in this example and the next one to avoid the confusion of having to double quote marks within the literals.) Embedded quotes and backslashes must be doubled:

```
SELECT $$the lexeme 'Joe"s' contains a quote$$::tsvector;
```
```
tsvectoR
---------------------------------------------------------------------
'Joe"s' 'a' 'contains' 'lexeme' 'quote' 'the'
```

Optionally, integer positions can be attached to lexemes:

```
SELECT 'a:1 fat:2 cat:3 sat:4 on:5 a:6 mat:7 and:8 ate:9 a:10 fat:11 rat:12'::tsvector;
```
Chapter 8. Data Types

A position normally indicates the source word’s location in the document. Positional information can be used for proximity ranking. Position values can range from 1 to 16383; larger numbers are silently set to 16383. Duplicate positions for the same lexeme are discarded.

Lexemes that have positions can further be labeled with a weight, which can be A, B, C, or D. D is the default and hence is not shown on output:

```
SELECT 'a:1A fat:2B,4C cat:5D'::tsvector;
```

Weights are typically used to reflect document structure, for example by marking title words differently from body words. Text search ranking functions can assign different priorities to the different weight markers.

It is important to understand that the tsvector type itself does not perform any word normalization; it assumes the words it is given are normalized appropriately for the application. For example,

```
SELECT 'The Fat Rats'::tsvector;
```

For most English-text-searching applications the above words would be considered non-normalized, but tsvector doesn’t care. Raw document text should usually be passed through to_tsvector to normalize the words appropriately for searching:

```
SELECT to_tsvector('english', 'The Fat Rats');
```

Again, see Chapter 12 for more detail.

8.11.2. tsquery

A tsquery value stores lexemes that are to be searched for, and can combine them using the Boolean operators & (AND), | (OR), and ! (NOT), as well as the phrase search operator <~> (FOLLOWED BY). There is also a variant <N> of the FOLLOWED BY operator, where N is an integer constant that specifies the distance between the two lexemes being searched for. <~> is equivalent to <1>.

Parentheses can be used to enforce grouping of these operators. In the absence of parentheses, ! (NOT) binds most tightly, <~> (FOLLOWED BY) next most tightly, then & (AND), with | (OR) binding the least tightly.

Here are some examples:

```
SELECT 'fat & rat'::tsquery;
```
SELECT 'fat & (rat | cat)':tsquery;
  tsquery
---------------------------
  'fat' & {'rat' | 'cat'}

SELECT 'fat & rat & ! cat':tsquery;
  tsquery
------------------------
  'fat' & 'rat' & !'cat'

Optionally, lexemes in a tsquery can be labeled with one or more weight letters, which restricts them to match only tsvector lexemes with one of those weights:

SELECT 'fat:ab & cat':tsquery;
  tsquery
------------------
  'fat':AB & 'cat'

Also, lexemes in a tsquery can be labeled with * to specify prefix matching:

SELECT 'super:*':tsquery;
  tsquery
-----------
  'super':*

This query will match any word in a tsvector that begins with "super".

Quoting rules for lexemes are the same as described previously for lexemes in tsvector; and, as with tsvector, any required normalization of words must be done before converting to the tsquery type. The to_tsquery function is convenient for performing such normalization:

SELECT to_tsquery('Fat:ab & Cats');
  to_tsquery
---------------
  'fat':AB & 'cat'

Note that to_tsquery will process prefixes in the same way as other words, which means this comparison returns true:

SELECT to_tsvector('postgraduate') @@ to_tsquery('postgres:*');
  to_tsvector | to_tsquery
---------------+------------
  'postgradu':1 | 'postgr':*

because postgres gets stemmed to postgr:

SELECT to_tsvector('postgraduate'), to_tsquery('postgres:*');
  to_tsvector | to_tsquery
---------------+------------
  'postgradu':1 | 'postgr':*

which will match the stemmed form of postgraduate.

150
8.12. UUID Type

The data type `uuid` stores Universally Unique Identifiers (UUID) as defined by RFC 4122, ISO/IEC 9834-8:2005, and related standards. (Some systems refer to this data type as a globally unique identifier, or GUID, instead.) This identifier is a 128-bit quantity that is generated by an algorithm chosen to make it very unlikely that the same identifier will be generated by anyone else in the known universe using the same algorithm. Therefore, for distributed systems, these identifiers provide a better uniqueness guarantee than sequence generators, which are only unique within a single database.

A UUID is written as a sequence of lower-case hexadecimal digits, in several groups separated by hyphens, specifically a group of 8 digits followed by three groups of 4 digits followed by a group of 12 digits, for a total of 32 digits representing the 128 bits. An example of a UUID in this standard form is:

```
a0eebc99-9c0b-4ef8-bb6d-6bb9bd380a11
```

PostgreSQL also accepts the following alternative forms for input: use of upper-case digits, the standard format surrounded by braces, omitting some or all hyphens, adding a hyphen after any group of four digits. Examples are:

```
A0EEBC99-9C0B-4EF8-BB6D-BB9BD380A11
{a0eebc99-9c0b-4ef8-bb6d-6bb9bd380a11}
a0ee-bc99-9c0b-4ef8-bb6d-6bb9-bd38-0a11
{a0eebc99-9c0b4ef8-bb6d6bb9-bd380a11}
```

Output is always in the standard form.

PostgreSQL provides storage and comparison functions for UUIDs, but the core database does not include any function for generating UUIDs, because no single algorithm is well suited for every application. The uuid-ossp module provides functions that implement several standard algorithms. The pgcrypto module also provides a generation function for random UUIDs. Alternatively, UUIDs could be generated by client applications or other libraries invoked through a server-side function.

8.13. XML Type

The `xml` data type can be used to store XML data. Its advantage over storing XML data in a text field is that it checks the input values for well-formedness, and there are support functions to perform type-safe operations on it; see Section 9.14. Use of this data type requires the installation to have been built with `configure --with-libxml`.

The `xml` type can store well-formed “documents”, as defined by the XML standard, as well as “content” fragments, which are defined by reference to the more permissive “document node”1 of the XQuery and XPath data model. Roughly, this means that content fragments can have more than one top-level element or character node. The expression `xmlvalue IS DOCUMENT` can be used to evaluate whether a particular `xml` value is a full document or only a content fragment.

8.13.1. Creating XML Values

To produce a value of type `xml` from character data, use the function `xmlparse`:

```
XMLPARSE ( { DOCUMENT | CONTENT } value)
```

Chapter 8. Data Types

Examples:

XMLPARSE (DOCUMENT '<?xml version="1.0"?><book><title>Manual</title><chapter>...</chapter></book>')
XMLPARSE (CONTENT 'abc<foo>bar</foo><bar>foo</bar>')

While this is the only way to convert character strings into XML values according to the SQL standard, the PostgreSQL-specific syntaxes:

xml '<foo>bar</foo>'
'<foo>bar</foo>'::xml

can also be used.

The xml type does not validate input values against a document type declaration (DTD), even when the input value specifies a DTD. There is also currently no built-in support for validating against other XML schema languages such as XML Schema.

The inverse operation, producing a character string value from xml, uses the function xmlserialize:

XMLSERIALIZE ( { DOCUMENT | CONTENT } value AS type )

type can be character, character varying, or text (or an alias for one of those). Again, according to the SQL standard, this is the only way to convert between type xml and character types, but PostgreSQL also allows you to simply cast the value.

When a character string value is cast to or from type xml without going through XMLPARSE or XMLSERIALIZE, respectively, the choice of DOCUMENT versus CONTENT is determined by the “XML option” session configuration parameter, which can be set using the standard command:

SET XML OPTION { DOCUMENT | CONTENT };

or the more PostgreSQL-like syntax

SET xmloption TO { DOCUMENT | CONTENT };

The default is CONTENT, so all forms of XML data are allowed.

8.13.2. Encoding Handling

Care must be taken when dealing with multiple character encodings on the client, server, and in the XML data passed through them. When using the text mode to pass queries to the server and query results to the client (which is the normal mode), PostgreSQL converts all character data passed between the client and the server and vice versa to the character encoding of the respective end; see Section 23.3. This includes string representations of XML values, such as in the above examples.

This would ordinarily mean that encoding declarations contained in XML data can become invalid as the character data is converted to other encodings while traveling between client and server, because the embedded encoding declaration is not changed. To cope with this behavior, encoding declarations contained in character strings presented for input to the xml type are ignored, and content is assumed to be in the current server encoding. Consequently, for correct processing, character strings of XML data must be sent from the client in the current client encoding. It is the responsibility of the client to either convert documents to the current client encoding before sending them to the server, or to adjust the client encoding appropriately. On output, values of type xml will not have an encoding declaration, and clients should assume all data is in the current client encoding.

When using binary mode to pass query parameters to the server and query results back to the client, no encoding conversion is performed, so the situation is different. In this case, an encoding declaration
in the XML data will be observed, and if it is absent, the data will be assumed to be in UTF-8 (as required by the XML standard; note that PostgreSQL does not support UTF-16). On output, data will have an encoding declaration specifying the client encoding, unless the client encoding is UTF-8, in which case it will be omitted.

Needless to say, processing XML data with PostgreSQL will be less error-prone and more efficient if the XML data encoding, client encoding, and server encoding are the same. Since XML data is internally processed in UTF-8, computations will be most efficient if the server encoding is also UTF-8.

**Caution**

Some XML-related functions may not work at all on non-ASCII data when the server encoding is not UTF-8. This is known to be an issue for `xpath()` in particular.

### 8.13.3. Accessing XML Values

The `xml` data type is unusual in that it does not provide any comparison operators. This is because there is no well-defined and universally useful comparison algorithm for XML data. One consequence of this is that you cannot retrieve rows by comparing an `xml` column against a search value. XML values should therefore typically be accompanied by a separate key field such as an ID. An alternative solution for comparing XML values is to convert them to character strings first, but note that character string comparison has little to do with a useful XML comparison method.

Since there are no comparison operators for the `xml` data type, it is not possible to create an index directly on a column of this type. If speedy searches in XML data are desired, possible workarounds include casting the expression to a character string type and indexing that, or indexing an XPath expression. Of course, the actual query would have to be adjusted to search by the indexed expression.

The text-search functionality in PostgreSQL can also be used to speed up full-document searches of XML data. The necessary preprocessing support is, however, not yet available in the PostgreSQL distribution.

### 8.14. JSON Types

JSON data types are for storing JSON (JavaScript Object Notation) data, as specified in RFC 7159\(^2\). Such data can also be stored as `text`, but the JSON data types have the advantage of enforcing that each stored value is valid according to the JSON rules. There are also assorted JSON-specific functions and operators available for data stored in these data types; see Section 9.15.

There are two JSON data types: `json` and `jsonb`. They accept almost identical sets of values as input. The major practical difference is one of efficiency. The `json` data type stores an exact copy of the input text, which processing functions must reparse on each execution; while `jsonb` data is stored in a decomposed binary format that makes it slightly slower to input due to added conversion overhead, but significantly faster to process, since no reparsing is needed. `jsonb` also supports indexing, which can be a significant advantage.

Because the `json` type stores an exact copy of the input text, it will preserve semantically-insignificant white space between tokens, as well as the order of keys within JSON objects. Also, if a JSON object

---

within the value contains the same key more than once, all the key/value pairs are kept. (The processing functions consider the last value as the operative one.) By contrast, \textit{jsonb} does not preserve white space, does not preserve the order of object keys, and does not keep duplicate object keys. If duplicate keys are specified in the input, only the last value is kept.

In general, most applications should prefer to store JSON data as \textit{jsonb}, unless there are quite specialized needs, such as legacy assumptions about ordering of object keys.

PostgreSQL allows only one character set encoding per database. It is therefore not possible for the \textit{JSON} types to conform rigidly to the JSON specification unless the database encoding is UTF8. Attempts to directly include characters that cannot be represented in the database encoding will fail; conversely, characters that can be represented in the database encoding but not in UTF8 will be allowed.

RFC 7159 permits JSON strings to contain Unicode escape sequences denoted by \texttt{\uXXXX}. In the input function for the \textit{json} type, Unicode escapes are allowed regardless of the database encoding, and are checked only for syntactic correctness (that is, that four hex digits follow \texttt{\u}). However, the input function for \textit{jsonb} is stricter: it disallows Unicode escapes for non-ASCII characters (those above \texttt{U+007F}) unless the database encoding is UTF8. The \textit{jsonb} type also rejects \texttt{\u0000} (because that cannot be represented in PostgreSQL's \textit{text} type), and it insists that any use of Unicode surrogate pairs to designate characters outside the Unicode Basic Multilingual Plane be correct. Valid Unicode escapes are converted to the equivalent ASCII or UTF8 character for storage; this includes folding surrogate pairs into a single character.

\textbf{Note:} Many of the JSON processing functions described in Section 9.15 will convert Unicode escapes to regular characters, and will therefore throw the same types of errors just described even if their input is of type \textit{json} not \textit{jsonb}. The fact that the \textit{json} input function does not make these checks may be considered a historical artifact, although it does allow for simple storage (without processing) of JSON Unicode escapes in a non-UTF8 database encoding. In general, it is best to avoid mixing Unicode escapes in JSON with a non-UTF8 database encoding, if possible.

When converting textual JSON input into \textit{jsonb}, the primitive types described by RFC 7159 are effectively mapped onto native PostgreSQL types, as shown in Table 8-23. Therefore, there are some minor additional constraints on what constitutes valid \textit{jsonb} data that do not apply to the \textit{json} type, nor to JSON in the abstract, corresponding to limits on what can be represented by the underlying data type. Notably, \textit{jsonb} will reject numbers that are outside the range of the PostgreSQL \textit{numeric} data type, while \textit{json} will not. Such implementation-defined restrictions are permitted by RFC 7159. However, in practice such problems are far more likely to occur in other implementations, as it is common to represent JSON’s \textit{number} primitive type as IEEE 754 double precision floating point (which RFC 7159 explicitly anticipates and allows for). When using JSON as an interchange format with such systems, the danger of losing numeric precision compared to data originally stored by PostgreSQL should be considered.

Conversely, as noted in the table there are some minor restrictions on the input format of JSON primitive types that do not apply to the corresponding PostgreSQL types.

\begin{table}[h]
\centering
\caption{JSON primitive types and corresponding PostgreSQL types}
\begin{tabular}{|l|l|l|}
\hline
\textbf{JSON primitive type} & \textbf{PostgreSQL type} & \textbf{Notes} \\
\hline
\texttt{string} & \texttt{text} & \texttt{\u0000} is disallowed, as are non-ASCII Unicode escapes if database encoding is not UTF8 \\
\hline
\end{tabular}
\end{table}
Chapter 8. Data Types

<table>
<thead>
<tr>
<th>JSON primitive type</th>
<th>PostgreSQL type</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>number</td>
<td>numeric</td>
<td>NaN and infinity values are disallowed</td>
</tr>
<tr>
<td>boolean</td>
<td>boolean</td>
<td>Only lowercase true and false spellings are accepted</td>
</tr>
<tr>
<td>null</td>
<td>(none)</td>
<td>SQL NULL is a different concept</td>
</tr>
</tbody>
</table>

8.14.1. JSON Input and Output Syntax

The input/output syntax for the JSON data types is as specified in RFC 7159.

The following are all valid json (or jsonb) expressions:

-- Simple scalar/primitive value
-- Primitive values can be numbers, quoted strings, true, false, or null
SELECT '5'::json;

-- Array of zero or more elements (elements need not be of same type)
SELECT '[1, 2, "foo", null]'::json;

-- Object containing pairs of keys and values
-- Note that object keys must always be quoted strings
SELECT '{"bar": "baz", "balance": 7.77, "active": false}'::json;

-- Arrays and objects can be nested arbitrarily
SELECT '{"foo": [true, "bar"], "tags": {"a": 1, "b": null}}'::json;

As previously stated, when a JSON value is input and then printed without any additional processing, json outputs the same text that was input, while jsonb does not preserve semantically-insignificant details such as whitespace. For example, note the differences here:

SELECT '{"bar": "baz", "balance": 7.77, "active":false}'::json;
json
-------------------------------------------------
{"bar": "baz", "balance": 7.77, "active":false}
(1 row)

SELECT '{"bar": "baz", "balance": 7.77, "active":false}'::jsonb;
jsonb
---------------------------------------------------
{"bar": "baz", "active": false, "balance": 7.77}
(1 row)

One semantically-insignificant detail worth noting is that in jsonb, numbers will be printed according to the behavior of the underlying numeric type. In practice this means that numbers entered with E notation will be printed without it, for example:

SELECT '{"reading": 1.230e-5}''::json, '{"reading": 1.230e-5}''::jsonb;
json          | jsonb
------------------------+------------------------
{"reading": 0.00001230} | {"reading": 0.00001230}
(1 row)
Chapter 8. Data Types

However, \texttt{jsonb} will preserve trailing fractional zeroes, as seen in this example, even though those are semantically insignificant for purposes such as equality checks.

8.14.2. Designing JSON documents effectively

Representing data as JSON can be considerably more flexible than the traditional relational data model, which is compelling in environments where requirements are fluid. It is quite possible for both approaches to co-exist and complement each other within the same application. However, even for applications where maximal flexibility is desired, it is still recommended that JSON documents have a somewhat fixed structure. The structure is typically unenforced (though enforcing some business rules declaratively is possible), but having a predictable structure makes it easier to write queries that usefully summarize a set of “documents” (datums) in a table.

JSON data is subject to the same concurrency-control considerations as any other data type when stored in a table. Although storing large documents is practicable, keep in mind that any update acquires a row-level lock on the whole row. Consider limiting JSON documents to a manageable size in order to decrease lock contention among updating transactions. Ideally, JSON documents should each represent an atomic datum that business rules dictate cannot reasonably be further subdivided into smaller datums that could be modified independently.

8.14.3. \texttt{jsonb} Containment and Existence

Testing containment is an important capability of \texttt{jsonb}. There is no parallel set of facilities for the \texttt{json} type. Containment tests whether one \texttt{jsonb} document has contained within it another one. These examples return true except as noted:

-- Simple scalar/primitive values contain only the identical value:
SELECT '"foo"::jsonb @> "foo"::jsonb;

-- The array on the right side is contained within the one on the left:
SELECT '[1, 2, 3]'::jsonb @> '[1, 3]'::jsonb;

-- Order of array elements is not significant, so this is also true:
SELECT '[1, 2, 3]'::jsonb @> '[3, 1]'::jsonb;

-- Duplicate array elements don’t matter either:
SELECT '[1, 2, 3]'::jsonb @> '[1, 2, 2]'::jsonb;

-- The object with a single pair on the right side is contained
-- within the object on the left side:
SELECT '{"product": "PostgreSQL", "version": 9.4, "jsonb": true}'::jsonb @> '{"version": 9.4}'::jsonb;

-- The array on the right side is \textbf{not} considered contained within the
-- array on the left, even though a similar array is nested within it:
SELECT '[1, 2, [1, 3]]'::jsonb @> '[1, 3]'::jsonb; -- yields false

-- But with a layer of nesting, it is contained:
SELECT '[1, 2, [1, 3]]'::jsonb @> '[[1, 3]]'::jsonb;

-- Similarly, containment is not reported here:
SELECT '{"foo": {"bar": "baz"}}'::jsonb @> '{"bar": "baz"}'::jsonb; -- yields false

-- A top-level key and an empty object is contained:
Chapter 8. Data Types
SELECT ’{"foo": {"bar": "baz"}}’::jsonb @> ’{"foo": {}}’::jsonb;

The general principle is that the contained object must match the containing object as to structure and
data contents, possibly after discarding some non-matching array elements or object key/value pairs
from the containing object. But remember that the order of array elements is not significant when
doing a containment match, and duplicate array elements are effectively considered only once.
As a special exception to the general principle that the structures must match, an array may contain a
primitive value:
-- This array contains the primitive string value:
SELECT ’["foo", "bar"]’::jsonb @> ’"bar"’::jsonb;
-- This exception is not reciprocal -- non-containment is reported here:
SELECT ’"bar"’::jsonb @> ’["bar"]’::jsonb; -- yields false
jsonb also has an existence operator, which is a variation on the theme of containment: it tests
whether a string (given as a text value) appears as an object key or array element at the top level of
the jsonb value. These examples return true except as noted:
-- String exists as array element:
SELECT ’["foo", "bar", "baz"]’::jsonb ? ’bar’;
-- String exists as object key:
SELECT ’{"foo": "bar"}’::jsonb ? ’foo’;
-- Object values are not considered:
SELECT ’{"foo": "bar"}’::jsonb ? ’bar’;

-- yields false

-- As with containment, existence must match at the top level:
SELECT ’{"foo": {"bar": "baz"}}’::jsonb ? ’bar’; -- yields false
-- A string is considered to exist if it matches a primitive JSON string:
SELECT ’"foo"’::jsonb ? ’foo’;

JSON objects are better suited than arrays for testing containment or existence when there are many
keys or elements involved, because unlike arrays they are internally optimized for searching, and do
not need to be searched linearly.
Tip: Because JSON containment is nested, an appropriate query can skip explicit selection of
sub-objects. As an example, suppose that we have a doc column containing objects at the top
level, with most objects containing tags fields that contain arrays of sub-objects. This query finds
entries in which sub-objects containing both "term":"paris" and "term":"food" appear, while
ignoring any such keys outside the tags array:
SELECT doc->’site_name’ FROM websites
WHERE doc @> ’{"tags":[{"term":"paris"}, {"term":"food"}]}’;

One could accomplish the same thing with, say,
SELECT doc->’site_name’ FROM websites
WHERE doc->’tags’ @> ’[{"term":"paris"}, {"term":"food"}]’;

but that approach is less flexible, and often less efficient as well.
On the other hand, the JSON existence operator is not nested: it will only look for the specified
key or array element at top level of the JSON value.

157


The various containment and existence operators, along with all other JSON operators and functions
are documented in Section 9.15.

8.14.4. jsonb Indexing

GIN indexes can be used to efficiently search for keys or key/value pairs occurring within a large
number of jsonb documents (datums). Two GIN “operator classes” are provided, offering different
performance and flexibility trade-offs.

The default GIN operator class for jsonb supports queries with top-level key-exists operators ?., ?% and
?| operators and path/value-exists operator @>. (For details of the semantics that these operators
implement, see Table 9-43.) An example of creating an index with this operator class is:

```
CREATE INDEX idxgin ON api USING GIN (jdoc);
```

The non-default GIN operator class jsonb_path_ops supports indexing the @> operator only. An
element of creating an index with this operator class is:

```
CREATE INDEX idxginp ON api USING GIN (jdoc jsonb_path_ops);
```

Consider the example of a table that stores JSON documents retrieved from a third-party web service,
with a documented schema definition. A typical document is:

```
{
    "guid": "9c36adc1-7fb5-4d5b-83b4-90356a46061a",
    "name": "Angela Barton",
    "is_active": true,
    "company": "Magnafone",
    "address": "178 Howard Place, Gulf, Washington, 702",
    "registered": "2009-11-07T08:53:22 +08:00",
    "latitude": 19.793713,
    "longitude": 86.513373,
    "tags": [
        "enim",
        "aliquip",
        "qui"
    ]
}
```

We store these documents in a table named api, in a jsonb column named jdoc. If a GIN index is
created on this column, queries like the following can make use of the index:

```
-- Find documents in which the key "company" has value "Magnafone"
SELECT jdoc->'guid', jdoc->'name' FROM api WHERE jdoc @> '{"company": "Magnafone"}';
```

However, the index could not be used for queries like the following, because though the operator ? is
indexable, it is not applied directly to the indexed column jdoc:

```
-- Find documents in which the key "tags" contains key or array element "qui"
SELECT jdoc->'guid', jdoc->'name' FROM api WHERE jdoc -> 'tags' ? 'qui';
```

Still, with appropriate use of expression indexes, the above query can use an index. If querying for
particular items within the "tags" key is common, defining an index like this may be worthwhile:

```
CREATE INDEX idxglntags ON api USING GIN ((jdoc -> 'tags'));
```
Now, the WHERE clause `jdoc -> 'tags' ? 'qui'` will be recognized as an application of the indexable operator `?` to the indexed expression `jdoc -> 'tags'`. (More information on expression indexes can be found in Section 11.7.)

Another approach to querying is to exploit containment, for example:

```
-- Find documents in which the key "tags" contains array element "qui"
SELECT jdoc->'guid', jdoc->'name' FROM api WHERE jdoc @> '{"tags": ["qui"]}';
```

A simple GIN index on the `jdoc` column can support this query. But note that such an index will store copies of every key and value in the `jdoc` column, whereas the expression index of the previous example stores only data found under the `tags` key. While the simple-index approach is far more flexible (since it supports queries about any key), targeted expression indexes are likely to be smaller and faster to search than a simple index.

Although the `jsonb_path_ops` operator class supports only queries with the `@>` operator, it has notable performance advantages over the default operator class `jsonb_ops`. A `jsonb_path_ops` index is usually much smaller than a `jsonb_ops` index over the same data, and the specificity of searches is better, particularly when queries contain keys that appear frequently in the data. Therefore search operations typically perform better than with the default operator class.

The technical difference between a `jsonb_ops` and a `jsonb_path_ops` GIN index is that the former creates independent index items for each key and value in the data, while the latter creates index items only for each value in the data. 1 Basically, each `jsonb_path_ops` index item is a hash of the value and the key(s) leading to it; for example to index `{"foo": {"bar": "baz"}}`, a single index item would be created incorporating all three of `foo`, `bar`, and `baz` into the hash value. Thus a containment query looking for this structure would result in an extremely specific index search; but there is no way at all to find out whether `foo` appears as a key. On the other hand, a `jsonb_ops` index would create three index items representing `foo`, `bar`, and `baz` separately; then to do the containment query, it would look for rows containing all three of these items. While GIN indexes can perform such an AND search fairly efficiently, it will still be less specific and slower than the equivalent `jsonb_path_ops` search, especially if there are a very large number of rows containing any single one of the three index items.

A disadvantage of the `jsonb_path_ops` approach is that it produces no index entries for JSON structures not containing any values, such as `{ "a": { } }`. If a search for documents containing such a structure is requested, it will require a full-index scan, which is quite slow. `jsonb_path_ops` is therefore ill-suited for applications that often perform such searches.

`jsonb` also supports `btree` and `hash` indexes. These are usually useful only if it’s important to check equality of complete JSON documents. The `btree` ordering for `jsonb` datums is seldom of great interest, but for completeness it is:

```
Object > Array > Boolean > Number > String > Null
Object with n pairs > object with n - 1 pairs
Array with n elements > array with n - 1 elements
```

Objects with equal numbers of pairs are compared in the order:

```
key-1, value-1, key-2 ...
```

Note that object keys are compared in their storage order; in particular, since shorter keys are stored before longer keys, this can lead to results that might be unintuitive, such as:

3. For this purpose, the term "value" includes array elements, though JSON terminology sometimes considers array elements distinct from values within objects.
Similarly, arrays with equal numbers of elements are compared in the order:

element-1, element-2 ...

Primitive JSON values are compared using the same comparison rules as for the underlying PostgreSQL data type. Strings are compared using the default database collation.

8.15. Arrays

PostgreSQL allows columns of a table to be defined as variable-length multidimensional arrays. Arrays of any built-in or user-defined base type, enum type, or composite type can be created. Arrays of domains are not yet supported.

8.15.1. Declaration of Array Types

To illustrate the use of array types, we create this table:

```sql
CREATE TABLE sal_emp (  
    name text,  
    pay_by_quarter integer[],  
    schedule text[][]
);
```

As shown, an array data type is named by appending square brackets ([ ]) to the data type name of the array elements. The above command will create a table named sal_emp with a column of type text (name), a one-dimensional array of type integer (pay_by_quarter), which represents the employee’s salary by quarter, and a two-dimensional array of text (schedule), which represents the employee’s weekly schedule.

The syntax for `CREATE TABLE` allows the exact size of arrays to be specified, for example:

```sql
CREATE TABLE tictactoe (  
    squares integer[3][3]
);
```

However, the current implementation ignores any supplied array size limits, i.e., the behavior is the same as for arrays of unspecified length.

The current implementation does not enforce the declared number of dimensions either. Arrays of a particular element type are all considered to be of the same type, regardless of size or number of dimensions. So, declaring the array size or number of dimensions in `CREATE TABLE` is simply documentation; it does not affect run-time behavior.

An alternative syntax, which conforms to the SQL standard by using the keyword `ARRAY`, can be used for one-dimensional arrays. `pay_by_quarter` could have been defined as:

```sql
pay_by_quarter integer ARRAY[4],
```

Or, if no array size is to be specified:

```sql
pay_by_quarter integer ARRAY,
```
Chapter 8. Data Types

As before, however, PostgreSQL does not enforce the size restriction in any case.

8.15.2. Array Value Input

To write an array value as a literal constant, enclose the element values within curly braces and separate them by commas. (If you know C, this is not unlike the C syntax for initializing structures.) You can put double quotes around any element value, and must do so if it contains commas or curly braces. (More details appear below.) Thus, the general format of an array constant is the following:

' { val1 delim val2 delim ... } '

where `delim` is the delimiter character for the type, as recorded in its `pg_type` entry. Among the standard data types provided in the PostgreSQL distribution, all use a comma (,), except for type `box` which uses a semicolon (;). Each `val` is either a constant of the array element type, or a subarray. An example of an array constant is:

' {1, 2, 3}, {4, 5, 6}, {7, 8, 9} '

This constant is a two-dimensional, 3-by-3 array consisting of three subarrays of integers.

To set an element of an array constant to NULL, write `NULL` for the element value. (Any upper- or lower-case variant of `NULL` will do.) If you want an actual string value "NULL", you must put double quotes around it.

(These kinds of array constants are actually only a special case of the generic type constants discussed in Section 4.1.2.7. The constant is initially treated as a string and passed to the array input conversion routine. An explicit type specification might be necessary.)

Now we can show some `INSERT` statements:

```sql
INSERT INTO sal_emp
VALUES ('Bill',
' {10000, 10000, 10000, 10000}',
' {{"meeting", "lunch"}, {{"training", "presentation"}}}');
```

```
INSERT INTO sal_emp
VALUES ('Carol',
' {20000, 25000, 25000, 25000}',
' {{"breakfast", "consulting"}, {{"meeting", "lunch"}}}');
```

The result of the previous two inserts looks like this:

```
SELECT * FROM sal_emp;
nname |  pay_by_quarter                  | schedule
-------+---------------------------+-------------------------------------------
    Bill | {10000, 10000, 10000, 10000} | {{meeting, lunch}, {training, presentation}}
   Carol | {20000, 25000, 25000, 25000} | {{breakfast, consulting}, {meeting, lunch}}
(2 rows)
```

Multidimensional arrays must have matching extents for each dimension. A mismatch causes an error, for example:

```sql
INSERT INTO sal_emp
VALUES ('Bill',
' {10000, 10000, 10000, 10000}',
' {{"meeting", "lunch"}, {{"training", "presentation"}}}');
```

```
ERROR:  invalid array subarray length
```
Chapter 8. Data Types

INSERT INTO sal_emp
VALUES ('Bill',
ARRAY[10000, 10000, 10000, 10000],
ARRAY[['meeting', 'lunch'], ['training', 'presentation']]));

INSERT INTO sal_emp
VALUES ('Carol',
ARRAY[20000, 25000, 25000, 25000],
ARRAY[['breakfast', 'consulting'], ['meeting', 'lunch']]));

Notice that the array elements are ordinary SQL constants or expressions; for instance, string literals are single quoted, instead of double quoted as they would be in an array literal. The ARRAY constructor syntax is discussed in more detail in Section 4.2.12.

8.15.3. Accessing Arrays

Now, we can run some queries on the table. First, we show how to access a single element of an array. This query retrieves the names of the employees whose pay changed in the second quarter:

SELECT name FROM sal_emp WHERE pay_by_quarter[1] <> pay_by_quarter[2];

name
-------
Carol
(1 row)

The array subscript numbers are written within square brackets. By default PostgreSQL uses a one-based numbering convention for arrays, that is, an array of n elements starts with array[1] and ends with array[n].

This query retrieves the third quarter pay of all employees:

SELECT pay_by_quarter[3] FROM sal_emp;

pay_by_quarter
-------------
10000
25000
(2 rows)

We can also access arbitrary rectangular slices of an array, or subarrays. An array slice is denoted by writing lower-bound:upper-bound for one or more array dimensions. For example, this query retrieves the first item on Bill’s schedule for the first two days of the week:

SELECT schedule[1:2][1:1] FROM sal_emp WHERE name = 'Bill';

schedule
Chapter 8. Data Types

---

({{meeting},{training}})
(1 row)

If any dimension is written as a slice, i.e., contains a colon, then all dimensions are treated as slices. Any dimension that has only a single number (no colon) is treated as being from 1 to the number specified. For example, [2] is treated as [1:2], as in this example:

SELECT schedule[1:2][2] FROM sal_emp WHERE name = 'Bill';

```
schedule
({{meeting,lunch},{training,presentation}})
(1 row)
```

To avoid confusion with the non-slice case, it’s best to use slice syntax for all dimensions, e.g., [1:2][1:1], not [2][1:1].

It is possible to omit the lower-bound and/or upper-bound of a slice specifier; the missing bound is replaced by the lower or upper limit of the array’s subscripts. For example:

SELECT schedule[:2][2:] FROM sal_emp WHERE name = 'Bill';

```
schedule
({{lunch},{presentation}})
(1 row)
```

SELECT schedule[:][1:1] FROM sal_emp WHERE name = 'Bill';

```
schedule
({{meeting},{training}})
(1 row)
```

An array subscript expression will return null if either the array itself or any of the subscript expressions are null. Also, null is returned if a subscript is outside the array bounds (this case does not raise an error). For example, if `schedule` currently has the dimensions [1:3][1:2] then referencing `schedule[3][3]` yields NULL. Similarly, an array reference with the wrong number of subscripts yields a null rather than an error.

An array slice expression likewise yields null if the array itself or any of the subscript expressions are null. However, in other cases such as selecting an array slice that is completely outside the current array bounds, a slice expression yields an empty (zero-dimensional) array instead of null. (This does not match non-slice behavior and is done for historical reasons.) If the requested slice partially overlaps the array bounds, then it is silently reduced to just the overlapping region instead of returning null.

The current dimensions of any array value can be retrieved with the `array_dims` function:

SELECT array_dims(schedule) FROM sal_emp WHERE name = 'Carol';

```
array_dims
-------------
[1:2][1:2]
(1 row)
```
array_dims produces a text result, which is convenient for people to read but perhaps inconvenient for programs. Dimensions can also be retrieved with `array_upper` and `array_lower`, which return the upper and lower bound of a specified array dimension, respectively:

```sql
SELECT array_upper(schedule, 1) FROM sal_emp WHERE name = 'Carol';
```

array_upper
-------------
2
(1 row)

array_length will return the length of a specified array dimension:

```sql
SELECT array_length(schedule, 1) FROM sal_emp WHERE name = 'Carol';
```

array_length
-------------
2
(1 row)

cardinality returns the total number of elements in an array across all dimensions. It is effectively the number of rows a call to `unnest` would yield:

```sql
SELECT cardinality(schedule) FROM sal_emp WHERE name = 'Carol';
```

cardinality
-------------
4
(1 row)

### 8.15.4. Modifying Arrays

An array value can be replaced completely:

```sql
UPDATE sal_emp SET pay_by_quarter = '{25000,25000,27000,27000}'
   WHERE name = 'Carol';
```

or using the `ARRAY` expression syntax:

```sql
UPDATE sal_emp SET pay_by_quarter = ARRAY[25000,25000,27000,27000]
   WHERE name = 'Carol';
```

An array can also be updated at a single element:

```sql
   WHERE name = 'Bill';
```

or updated in a slice:

```sql
UPDATE sal_emp SET pay_by_quarter[1:2] = '{27000,27000}'
   WHERE name = 'Carol';
```
Chapter 8. Data Types

The slice syntaxes with omitted lower-bound and/or upper-bound can be used too, but only when updating an array value that is not NULL or zero-dimensional (otherwise, there is no existing subscript limit to substitute).

A stored array value can be enlarged by assigning to elements not already present. Any positions between those previously present and the newly assigned elements will be filled with nulls. For example, if array myarray currently has 4 elements, it will have six elements after an update that assigns to myarray[6]; myarray[5] will contain null. Currently, enlargement in this fashion is only allowed for one-dimensional arrays, not multidimensional arrays.

Subscripted assignment allows creation of arrays that do not use one-based subscripts. For example one might assign to myarray[-2:7] to create an array with subscript values from -2 to 7.

New array values can also be constructed using the concatenation operator, ||:

```sql
SELECT ARRAY[1,2] || ARRAY[3,4];
?column?
-------------------
(1,2,3,4)  (1 row)
```

```sql
SELECT ARRAY[5,6] || ARRAY[[1,2],[3,4]];
?column?
---------------------
(5,6),(1,2),(3,4)  (1 row)
```

The concatenation operator allows a single element to be pushed onto the beginning or end of a one-dimensional array. It also accepts two N-dimensional arrays, or an N-dimensional and an N+1-dimensional array.

When a single element is pushed onto either the beginning or end of a one-dimensional array, the result is an array with the same lower bound subscript as the array operand. For example:

```sql
SELECT array_dims(1 || ’[0:1]’={2,3}::int[]);
array_dims
------------
[0:2]  (1 row)
```

```sql
SELECT array_dims(ARRAY[1,2] || 3);
array_dims
------------
[1:3]  (1 row)
```

When two arrays with an equal number of dimensions are concatenated, the result retains the lower bound subscript of the left-hand operand’s outer dimension. The result is an array comprising every element of the left-hand operand followed by every element of the right-hand operand. For example:

```sql
SELECT array_dims(ARRAY[1,2] || ARRAY[3,4,5]);
array_dims
------------
[1:5]  (1 row)
```
Chapter 8. Data Types

```
SELECT array_dims(ARRAY[[1,2],[3,4]] || ARRAY[[5,6],[7,8],[9,0]]);
array_dims
------------
[1:5][1:2]
(1 row)

When an $N$-dimensional array is pushed onto the beginning or end of an $N+1$-dimensional array, the result is analogous to the element-array case above. Each $N$-dimensional sub-array is essentially an element of the $N+1$-dimensional array’s outer dimension. For example:

```
SELECT array_dims(ARRAY[1,2] || ARRAY[[3,4],[5,6]]);
array_dims
------------
[1:3][1:2]
(1 row)
```

An array can also be constructed by using the functions `array_prepend`, `array_append`, or `array_cat`. The first two only support one-dimensional arrays, but `array_cat` supports multidimensional arrays. Some examples:

```
SELECT array_prepend(1, ARRAY[2,3]);
array_prepend
---------------
{1,2,3}
(1 row)

SELECT array_append(ARRAY[1,2], 3);
array_append
------------
{1,2,3}
(1 row)

SELECT array_cat(ARRAY[1,2], ARRAY[3,4]);
array_cat
----------
{1,2,3,4}
(1 row)

SELECT array_cat(ARRAY[[1,2],[3,4]], ARRAY[5,6]);
array_cat
-------------
{{{1,2},{3,4},{5,6}}
(1 row)

SELECT array_cat(ARRAY[5,6], ARRAY[[1,2],[3,4]]);
array_cat
--------------
{{5,6},{1,2},{3,4}}
```
In simple cases, the concatenation operator discussed above is preferred over direct use of these functions. However, because the concatenation operator is overloaded to serve all three cases, there are situations where use of one of the functions is helpful to avoid ambiguity. For example consider:

```sql
SELECT ARRAY[1, 2] || '{3, 4}'; -- the untyped literal is taken as an array
---------
{1,2,3,4}

SELECT ARRAY[1, 2] || '7'; -- so is this one
ERROR: malformed array literal: "7"

SELECT ARRAY[1, 2] || NULL; -- so is an undecorated NULL
---------
{1,2}
(1 row)

SELECT array_append(ARRAY[1, 2], NULL); -- this might have been meant
array_append
---------
{1,2,NULL}
```

In the examples above, the parser sees an integer array on one side of the concatenation operator, and a constant of undetermined type on the other. The heuristic it uses to resolve the constant’s type is to assume it’s of the same type as the operator’s other input — in this case, integer array. So the concatenation operator is presumed to represent `array_cat`, not `array_append`. When that’s the wrong choice, it could be fixed by casting the constant to the array’s element type; but explicit use of `array_append` might be a preferable solution.

### 8.15.5. Searching in Arrays

To search for a value in an array, each value must be checked. This can be done manually, if you know the size of the array. For example:

```sql
SELECT * FROM sal_emp WHERE pay_by_quarter[1] = 10000 OR
  pay_by_quarter[2] = 10000 OR
  pay_by_quarter[3] = 10000 OR
  pay_by_quarter[4] = 10000;
```

However, this quickly becomes tedious for large arrays, and is not helpful if the size of the array is unknown. An alternative method is described in Section 9.23. The above query could be replaced by:

```sql
SELECT * FROM sal_emp WHERE 10000 = ANY (pay_by_quarter);
```

In addition, you can find rows where the array has all values equal to 10000 with:

```sql
SELECT * FROM sal_emp WHERE 10000 = ALL (pay_by_quarter);
```

Alternatively, the `generate_subscripts` function can be used. For example:

```sql
SELECT * FROM
  (SELECT pay_by_quarter,
       generate_subscripts(pay_by_quarter, 1) AS s
```

167
FROM sal_emp) AS foo
WHERE pay_by_quarter[s] = 10000;

This function is described in Table 9-58.

You can also search an array using the \&\& operator, which checks whether the left operand overlaps with the right operand. For instance:

\[\text{SELECT } \ast \text{ FROM sal_emp WHERE pay_by_quarter } \&\& \text{ ARRAY[10000];}\]

This and other array operators are further described in Section 9.18. It can be accelerated by an appropriate index, as described in Section 11.2.

You can also search for specific values in an array using the array_position and array_positions functions. The former returns the subscript of the first occurrence of a value in an array; the latter returns an array with the subscripts of all occurrences of the value in the array. For example:

\[\text{SELECT array_position(ARRAY[\text{\textquote{sun}},\text{\textquote{mon}},\text{\textquote{tue}},\text{\textquote{wed}},\text{\textquote{thu}},\text{\textquote{fri}},\text{\textquote{sat}}], \text{\textquote{mon}}); array_positions}
\]

\[\text{-----------}
\]

\[2\]

\[\text{SELECT array_positions(ARRAY[1, 4, 3, 1, 3, 4, 2, 1], 1);}\]

\[\text{array_positions}
\]

\[\text{-----------------}
\]

\[\{1, 4, 8\}\]

Tip: Arrays are not sets; searching for specific array elements can be a sign of database misdesign. Consider using a separate table with a row for each item that would be an array element. This will be easier to search, and is likely to scale better for a large number of elements.

8.15.6. Array Input and Output Syntax

The external text representation of an array value consists of items that are interpreted according to the I/O conversion rules for the array’s element type, plus decoration that indicates the array structure. The decoration consists of curly braces (\{ and \}) around the array value plus delimiter characters between adjacent items. The delimiter character is usually a comma (, ) but can be something else: it is determined by the typdelim setting for the array’s element type. Among the standard data types provided in the PostgreSQL distribution, all use a comma, except for type box, which uses a semicolon (;). In a multidimensional array, each dimension (row, plane, cube, etc.) gets its own level of curly braces, and delimiters must be written between adjacent curly-braced entities of the same level.

The array output routine will put double quotes around element values if they are empty strings, contain curly braces, delimiter characters, double quotes, backslashes, or white space, or match the word NULL. Double quotes and backslashes embedded in element values will be backslash-escaped. For numeric data types it is safe to assume that double quotes will never appear, but for textual data types one should be prepared to cope with either the presence or absence of quotes.

By default, the lower bound index value of an array’s dimensions is set to one. To represent arrays with other lower bounds, the array subscript ranges can be specified explicitly before writing the array.
This decoration consists of square brackets ([]) around each array dimension’s lower and upper bounds, with a colon (:) delimiter character in between. The array dimension decoration is followed by an equal sign (=). For example:

```
SELECT f1[1][-2][3] AS e1, f1[1][-1][5] AS e2
FROM (SELECT '[1:1][-2:-1][3:5]={{{1,2,3},{4,5,6}}}'::int[] AS f1) AS ss;
```

<table>
<thead>
<tr>
<th>e1</th>
<th>e2</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>6</td>
</tr>
</tbody>
</table>

(1 row)

The array output routine will include explicit dimensions in its result only when there are one or more lower bounds different from one.

If the value written for an element is `NULL` (in any case variant), the element is taken to be `NULL`. The presence of any quotes or backslashes disables this and allows the literal string value “NULL” to be entered. Also, for backward compatibility with pre-8.2 versions of PostgreSQL, the `array_nulls` configuration parameter can be turned off to suppress recognition of `NULL` as a `NULL`.

As shown previously, when writing an array value you can use double quotes around any individual array element. You must do so if the element value would otherwise confuse the array-value parser. For example, elements containing curly braces, commas (or the data type’s delimiter character), double quotes, backslashes, or leading or trailing whitespace must be double-quoted. Empty strings and strings matching the word `NULL` must be quoted, too. To put a double quote or backslash in a quoted array element value, precede it with a backslash. Alternatively, you can avoid quotes and use backslash-escaping to protect all data characters that would otherwise be taken as array syntax.

You can add whitespace before a left brace or after a right brace. You can also add whitespace before or after any individual item string. In all of these cases the whitespace will be ignored. However, whitespace within double-quoted elements, or surrounded on both sides by non-whitespace characters of an element, is not ignored.

**Tip:** The `ARRAY` constructor syntax (see Section 4.2.12) is often easier to work with than the array-literal syntax when writing array values in SQL commands. In `ARRAY`, individual element values are written the same way they would be written when not members of an array.

### 8.16. Composite Types

A **composite type** represents the structure of a row or record; it is essentially just a list of field names and their data types. PostgreSQL allows composite types to be used in many of the same ways that simple types can be used. For example, a column of a table can be declared to be of a composite type.

#### 8.16.1. Declaration of Composite Types

Here are two simple examples of defining composite types:

```
CREATE TYPE complex AS (
    r double precision,
    i double precision
);
```
CREATE TYPE inventory_item AS (
    name text,
    supplier_id integer,
    price numeric
);

The syntax is comparable to CREATE TABLE, except that only field names and types can be specified; no constraints (such as NOT NULL) can presently be included. Note that the AS keyword is essential; without it, the system will think a different kind of CREATE TYPE command is meant, and you will get odd syntax errors.

Having defined the types, we can use them to create tables:

CREATE TABLE on_hand (  
    item inventory_item,
    count integer
);

INSERT INTO on_hand VALUES (ROW('fuzzy dice', 42, 1.99), 1000);

or functions:

CREATE FUNCTION price_extension(inventory_item, integer) RETURNS numeric  
AS 'SELECT $1.price * $2' LANGUAGE SQL;

SELECT price_extension(item, 10) FROM on_hand;

Whenever you create a table, a composite type is also automatically created, with the same name as the table, to represent the table’s row type. For example, had we said:

CREATE TABLE inventory_item (  
    name text,
    supplier_id integer REFERENCES suppliers,
    price numeric CHECK (price > 0)
);

then the same inventory_item composite type shown above would come into being as a byproduct, and could be used just as above. Note however an important restriction of the current implementation: since no constraints are associated with a composite type, the constraints shown in the table definition do not apply to values of the composite type outside the table. (A partial workaround is to use domain types as members of composite types.)

8.16.2. Constructing Composite Values

To write a composite value as a literal constant, enclose the field values within parentheses and separate them by commas. You can put double quotes around any field value, and must do so if it contains commas or parentheses. (More details appear below.) Thus, the general format of a composite constant is the following:

'( val1, val2, ... )'

An example is:
Chapter 8. Data Types

'("fuzzy dice",42,1.99)'

which would be a valid value of the inventory_item type defined above. To make a field be NULL, write no characters at all in its position in the list. For example, this constant specifies a NULL third field:

'("fuzzy dice",42, )'

If you want an empty string rather than NULL, write double quotes:

'("",42, )'

Here the first field is a non-NULL empty string, the third is NULL.

(These constants are actually only a special case of the generic type constants discussed in Section 4.1.2.7. The constant is initially treated as a string and passed to the composite-type input conversion routine. An explicit type specification might be necessary to tell which type to convert the constant to.)

The ROW expression syntax can also be used to construct composite values. In most cases this is considerably simpler to use than the string-literal syntax since you don’t have to worry about multiple layers of quoting. We already used this method above:

ROW('fuzzy dice‘, 42, 1.99)
ROW("", 42, NULL)

The ROW keyword is actually optional as long as you have more than one field in the expression, so these can be simplified to:

('fuzzy dice‘, 42, 1.99)
("", 42, NULL)

The ROW expression syntax is discussed in more detail in Section 4.2.13.

8.16.3. Accessing Composite Types

To access a field of a composite column, one writes a dot and the field name, much like selecting a field from a table name. In fact, it’s so much like selecting from a table name that you often have to use parentheses to keep from confusing the parser. For example, you might try to select some subfields from our on_hand example table with something like:

SELECT item.name FROM on_hand WHERE item.price > 9.99;

This will not work since the name item is taken to be a table name, not a column name of on_hand, per SQL syntax rules. You must write it like this:

SELECT (item).name FROM on_hand WHERE (item).price > 9.99;

or if you need to use the table name as well (for instance in a multitable query), like this:

SELECT (on_hand.item).name FROM on_hand WHERE (on_hand.item).price > 9.99;

Now the parenthesized object is correctly interpreted as a reference to the item column, and then the subfield can be selected from it.
Similar syntactic issues apply whenever you select a field from a composite value. For instance, to select just one field from the result of a function that returns a composite value, you’d need to write something like:

```
SELECT (my_func(...)).field FROM ...
```

Without the extra parentheses, this will generate a syntax error.

The special field name * means “all fields”, as further explained in Section 8.16.5.

### 8.16.4. Modifying Composite Types

Here are some examples of the proper syntax for inserting and updating composite columns. First, inserting or updating a whole column:

```
INSERT INTO mytab (complex_col) VALUES((1.1,2.2));
```

```
UPDATE mytab SET complex_col = ROW(1.1,2.2) WHERE ...;
```

The first example omits `ROW`, the second uses it; we could have done it either way.

We can update an individual subfield of a composite column:

```
UPDATE mytab SET complex_col.r = (complex_col).r + 1 WHERE ...;
```

Notice here that we don’t need to (and indeed cannot) put parentheses around the column name appearing just after `SET`, but we do need parentheses when referencing the same column in the expression to the right of the equal sign.

And we can specify subfields as targets for `INSERT`, too:

```
INSERT INTO mytab (complex_col.r, complex_col.i) VALUES(1.1, 2.2);
```

Had we not supplied values for all the subfields of the column, the remaining subfields would have been filled with null values.

### 8.16.5. Using Composite Types in Queries

There are various special syntax rules and behaviors associated with composite types in queries. These rules provide useful shortcuts, but can be confusing if you don’t know the logic behind them.

In PostgreSQL, a reference to a table name (or alias) in a query is effectively a reference to the composite value of the table’s current row. For example, if we had a table `inventory_item` as shown above, we could write:

```
SELECT c FROM inventory_item c;
```

This query produces a single composite-valued column, so we might get output like:

```
c---------------------------------
  ("fuzzy dice",42,1.99)
(1 row)
```

Note however that simple names are matched to column names before table names, so this example works only because there is no column named `c` in the query’s tables.
The ordinary qualified-column-name syntax `table_name.column_name` can be understood as applying field selection to the composite value of the table’s current row. (For efficiency reasons, it’s not actually implemented that way.)

When we write

```sql
SELECT c.* FROM inventory_item c;
```

then, according to the SQL standard, we should get the contents of the table expanded into separate columns:

<table>
<thead>
<tr>
<th>name</th>
<th>supplier_id</th>
<th>price</th>
</tr>
</thead>
<tbody>
<tr>
<td>fuzzy dice</td>
<td>42</td>
<td>1.99</td>
</tr>
</tbody>
</table>

(1 row)

as if the query were

```sql
SELECT c.name, c.supplier_id, c.price FROM inventory_item c;
```

PostgreSQL will apply this expansion behavior to any composite-valued expression, although as shown above, you need to write parentheses around the value that `*` is applied to whenever it’s not a simple table name. For example, if `myfunc()` is a function returning a composite type with columns `a`, `b`, and `c`, then these two queries have the same result:

```sql
SELECT (myfunc(x)).* FROM some_table;
SELECT (myfunc(x)).a, (myfunc(x)).b, (myfunc(x)).c FROM some_table;
```

**Tip:** PostgreSQL handles column expansion by actually transforming the first form into the second. So, in this example, `myfunc()` would get invoked three times per row with either syntax. If it’s an expensive function you may wish to avoid that, which you can do with a query like:

```sql
SELECT (m).* FROM (SELECT myfunc(x) AS m FROM some_table OFFSET 0) ss;
```

The `OFFSET 0` clause keeps the optimizer from “flattening” the sub-select to arrive at the form with multiple calls of `myfunc()`.

The `composite_value.*` syntax results in column expansion of this kind when it appears at the top level of a `SELECT` output list, a `RETURNING` list in `INSERT/UPDATE/DELETE`, a `VALUES` clause, or a row constructor. In all other contexts (including when nested inside one of those constructs), attaching `.*` to a composite value does not change the value, since it means “all columns” and so the same composite value is produced again. For example, if `somefunc()` accepts a composite-valued argument, these queries are the same:

```sql
SELECT somefunc(c.*) FROM inventory_item c;
SELECT somefunc(c) FROM inventory_item c;
```

In both cases, the current row of `inventory_item` is passed to the function as a single composite-valued argument. Even though `.*` does nothing in such cases, using it is good style, since it makes clear that a composite value is intended. In particular, the parser will consider `c` in `c.*` to refer to a table name or alias, not to a column name, so that there is no ambiguity; whereas without `.*`, it is not clear whether `c` means a table name or a column name, and in fact the column-name interpretation will be preferred if there is a column named `c`.  

173
Another example demonstrating these concepts is that all these queries mean the same thing:

```
SELECT * FROM inventory_item c ORDER BY c;
SELECT * FROM inventory_item c ORDER BY c.*;
SELECT * FROM inventory_item c ORDER BY ROW(c.*);
```

All of these ORDER BY clauses specify the row’s composite value, resulting in sorting the rows according to the rules described in Section 9.23.6. However, if `inventory_item` contained a column named `c`, the first case would be different from the others, as it would mean to sort by that column only. Given the column names previously shown, these queries are also equivalent to those above:

```
SELECT * FROM inventory_item c ORDER BY ROW(c.name, c.supplier_id, c.price);
SELECT * FROM inventory_item c ORDER BY (c.name, c.supplier_id, c.price);
```

(The last case uses a row constructor with the key word `ROW` omitted.)

Another special syntactical behavior associated with composite values is that we can use functional notation for extracting a field of a composite value. The simple way to explain this is that the notations `field(table)` and `table.field` are interchangeable. For example, these queries are equivalent:

```
SELECT c.name FROM inventory_item c WHERE c.price > 1000;
SELECT name(c) FROM inventory_item c WHERE price(c) > 1000;
```

Moreover, if we have a function that accepts a single argument of a composite type, we can call it with either notation. These queries are all equivalent:

```
SELECT somefunc(c) FROM inventory_item c;
SELECT somefunc(c.*) FROM inventory_item c;
SELECT c.somefunc FROM inventory_item c;
```

This equivalence between functional notation and field notation makes it possible to use functions on composite types to implement “computed fields”. An application using the last query above wouldn’t need to be directly aware that `somefunc` isn’t a real column of the table.

**Tip:** Because of this behavior, it’s unwise to give a function that takes a single composite-type argument the same name as any of the fields of that composite type. If there is ambiguity, the field-name interpretation will be preferred, so that such a function could not be called without tricks. One way to force the function interpretation is to schema-qualify the function name, that is, write `schema.func(compositevalue)`.

### 8.16.6. Composite Type Input and Output Syntax

The external text representation of a composite value consists of items that are interpreted according to the I/O conversion rules for the individual field types, plus decoration that indicates the composite structure. The decoration consists of parentheses (`(` and `)`) around the whole value, plus commas (`,`) between adjacent items. Whitespace outside the parentheses is ignored, but within the parentheses it is considered part of the field value, and might or might not be significant depending on the input conversion rules for the field data type. For example, in:

```
'( 42)'
```

the whitespace will be ignored if the field type is integer, but not if it is text.
As shown previously, when writing a composite value you can write double quotes around any individual field value. You must do so if the field value would otherwise confuse the composite-value parser. In particular, fields containing parentheses, commas, double quotes, or backslashes must be double-quoted. To put a double quote or backslash in a quoted composite field value, precede it with a backslash. (Also, a pair of double quotes within a double-quoted field value is taken to represent a double quote character, analogously to the rules for single quotes in SQL literal strings.) Alternatively, you can avoid quoting and use backslash-escaping to protect all data characters that would otherwise be taken as composite syntax.

A completely empty field value (no characters at all between the commas or parentheses) represents a NULL. To write a value that is an empty string rather than NULL, write "".

The composite output routine will put double quotes around field values if they are empty strings or contain parentheses, commas, double quotes, backslashes, or white space. (Doing so for white space is not essential, but aids legibility.) Double quotes and backslashes embedded in field values will be doubled.

**Note:** Remember that what you write in an SQL command will first be interpreted as a string literal, and then as a composite. This doubles the number of backslashes you need (assuming escape string syntax is used). For example, to insert a text field containing a double quote and a backslash in a composite value, you'd need to write:

```sql
INSERT ... VALUES ('("\")');
```

The string-literal processor removes one level of backslashes, so that what arrives at the composite-value parser looks like "("\")". In turn, the string fed to the text data type’s input routine becomes ". (If we were working with a data type whose input routine also treated backslashes specially, for example, we might need as many as eight backslashes in the command to get one backslash into the stored composite field.) Dollar quoting (see Section 4.1.2.4) can be used to avoid the need to double backslashes.

**Tip:** The `ROW` constructor syntax is usually easier to work with than the composite-literal syntax when writing composite values in SQL commands. In `ROW`, individual field values are written the same way they would be written when not members of a composite.

### 8.17. Range Types

Range types are data types representing a range of values of some element type (called the range’s subtype). For instance, ranges of `timestamp` might be used to represent the ranges of time that a meeting room is reserved. In this case the data type is `tsrange` (short for “timestamp range”), and `timestamp` is the subtype. The subtype must have a total order so that it is well-defined whether element values are within, before, or after a range of values.

Range types are useful because they represent many element values in a single range value, and because concepts such as overlapping ranges can be expressed clearly. The use of time and date ranges for scheduling purposes is the clearest example; but price ranges, measurement ranges from an instrument, and so forth can also be useful.
Chapter 8. Data Types

8.17.1. Built-in Range Types

PostgreSQL comes with the following built-in range types:

- `int4range` — Range of integer
- `int8range` — Range of bigint
- `numrange` — Range of numeric
- `tsrange` — Range of timestamp without time zone
- `tstzrange` — Range of timestamp with time zone
- `daterange` — Range of date

In addition, you can define your own range types; see CREATE TYPE for more information.

8.17.2. Examples

CREATE TABLE reservation (room int, during tsrange);
INSERT INTO reservation VALUES
(1108, '2010-01-01 14:30, 2010-01-01 15:30');

-- Containment
SELECT int4range(10, 20) @> 3;

-- Overlaps
SELECT numrange(11.1, 22.2) && numrange(20.0, 30.0);

-- Extract the upper bound
SELECT upper(int8range(15, 25));

-- Compute the intersection
SELECT int4range(10, 20) * int4range(15, 25);

-- Is the range empty?
SELECT isempty(numrange(1, 5));

See Table 9-49 and Table 9-50 for complete lists of operators and functions on range types.

8.17.3. Inclusive and Exclusive Bounds

Every non-empty range has two bounds, the lower bound and the upper bound. All points between these values are included in the range. An inclusive bound means that the boundary point itself is included in the range as well, while an exclusive bound means that the boundary point is not included in the range.

In the text form of a range, an inclusive lower bound is represented by `"["` while an exclusive lower bound is represented by `"("`. Likewise, an inclusive upper bound is represented by `"]"`, while an exclusive upper bound is represented by `"\)"`. (See Section 8.17.5 for more details.)

The functions `lower_inc` and `upper_inc` test the inclusivity of the lower and upper bounds of a range value, respectively.
8.17.4. Infinite (Unbounded) Ranges

The lower bound of a range can be omitted, meaning that all values less than the upper bound are included in the range, e.g., \((, 3]\). Likewise, if the upper bound of the range is omitted, then all values greater than the lower bound are included in the range. If both lower and upper bounds are omitted, all values of the element type are considered to be in the range. Specifying a missing bound as inclusive is automatically converted to exclusive, e.g., \([, ]\) is converted to \((, )\). You can think of these missing values as $\pm\infty$, but they are special range type values and are considered to be beyond any range element type’s $\pm\infty$ values.

Element types that have the notion of “infinity” can use them as explicit bound values. For example, with timestamp ranges, \([\text{today}, \infty)\) excludes the special timestamp value \(\infty\), while \([\text{today}, \infty]\) includes it, as does \([\text{today}, )\) and \([\text{today}, ]\).

The functions \(\text{lower}\_\text{inf}\) and \(\text{upper}\_\text{inf}\) test for infinite lower and upper bounds of a range, respectively.

8.17.5. Range Input/Output

The input for a range value must follow one of the following patterns:

\[
\begin{align*}
& (\text{lower-bound}, \text{upper-bound}) \\
& (\text{lower-bound}, \text{upper-bound}] \\
& [\text{lower-bound}, \text{upper-bound}) \\
& [\text{lower-bound}, \text{upper-bound}] \\
& \text{empty}
\end{align*}
\]

The parentheses or brackets indicate whether the lower and upper bounds are exclusive or inclusive, as described previously. Notice that the final pattern is \(\text{empty}\), which represents an empty range (a range that contains no points).

The \(\text{lower-bound}\) may be either a string that is valid input for the subtype, or empty to indicate no lower bound. Likewise, \(\text{upper-bound}\) may be either a string that is valid input for the subtype, or empty to indicate no upper bound.

Each bound value can be quoted using " (double quote) characters. This is necessary if the bound value contains parentheses, brackets, commas, double quotes, or backslashes, since these characters would otherwise be taken as part of the range syntax. To put a double quote or backslash in a quoted bound value, precede it with a backslash. (Also, a pair of double quotes within a double-quoted bound value is taken to represent a double quote character, analogously to the rules for single quotes in SQL literal strings.) Alternatively, you can avoid quoting and use backslash-escaping to protect all data characters that would otherwise be taken as range syntax. Also, to write a bound value that is an empty string, write "", since writing nothing means an infinite bound.

Whitespace is allowed before and after the range value, but any whitespace between the parentheses or brackets is taken as part of the lower or upper bound value. (Depending on the element type, it might or might not be significant.)

Note: These rules are very similar to those for writing field values in composite-type literals. See Section 8.16.6 for additional commentary.

Examples:

\(--\) includes 3, does not include 7, and does include all points in between
Chapter 8. Data Types

SELECT '3,7'::int4range;
-- does not include either 3 or 7, but includes all points in between
SELECT '(3,7)'::int4range;
-- includes only the single point 4
SELECT '{4,4}'::int4range;
-- includes no points (and will be normalized to 'empty')
SELECT '{4,4}'::int4range;

8.17.6. Constructing Ranges

Each range type has a constructor function with the same name as the range type. Using the constructor function is frequently more convenient than writing a range literal constant, since it avoids the need for extra quoting of the bound values. The constructor function accepts two or three arguments. The two-argument form constructs a range in standard form (lower bound inclusive, upper bound exclusive), while the three-argument form constructs a range with bounds of the form specified by the third argument. The third argument must be one of the strings "()", "["", "]", or "[]". For example:

-- The full form is: lower bound, upper bound, and text argument indicating
-- inclusivity/exclusivity of bounds.
SELECT numrange(1.0, 14.0, '()');

-- If the third argument is omitted, '()' is assumed.
SELECT numrange(1.0, 14.0);

-- Although '() is specified here, on display the value will be converted to
-- canonical form, since int8range is a discrete range type (see below).
SELECT int8range(1, 14, '()');

-- Using NULL for either bound causes the range to be unbounded on that side.
SELECT numrange(NULL, 2.2);

8.17.7. Discrete Range Types

A discrete range is one whose element type has a well-defined “step”, such as integer or date. In these types two elements can be said to be adjacent, when there are no valid values between them. This contrasts with continuous ranges, where it’s always (or almost always) possible to identify other element values between two given values. For example, a range over the numeric type is continuous, as is a range over timestamp. (Even though timestamp has limited precision, and so could theoretically be treated as discrete, it’s better to consider it continuous since the step size is normally not of interest.)

Another way to think about a discrete range type is that there is a clear idea of a “next” or “previous” value for each element value. Knowing that, it is possible to convert between inclusive and exclusive representations of a range’s bounds, by choosing the next or previous element value instead of the
one originally given. For example, in an integer range type \([4, 8]\) and \((3, 9)\) denote the same set of values; but this would not be so for a range over numeric.

A discrete range type should have a canonicalization function that is aware of the desired step size for the element type. The canonicalization function is charged with converting equivalent values of the range type to have identical representations, in particular consistently inclusive or exclusive bounds. If a canonicalization function is not specified, then ranges with different formatting will always be treated as unequal, even though they might represent the same set of values in reality.

The built-in range types int4range, int8range, and daterange all use a canonical form that includes the lower bound and excludes the upper bound; that is, \((,)\). User-defined range types can use other conventions, however.

### 8.17.8. Defining New Range Types

Users can define their own range types. The most common reason to do this is to use ranges over subtypes not provided among the built-in range types. For example, to define a new range type of subtype float8:

```
CREATE TYPE floatrange AS RANGE (
  subtype = float8,
  subtype_diff = float8mi
);
```

```
SELECT '([1.234, 5.678]::floatrange;
```

Because float8 has no meaningful “step”, we do not define a canonicalization function in this example.

Defining your own range type also allows you to specify a different subtype B-tree operator class or collation to use, so as to change the sort ordering that determines which values fall into a given range.

If the subtype is considered to have discrete rather than continuous values, the `CREATE TYPE` command should specify a canonical function. The canonicalization function takes an input range value, and must return an equivalent range value that may have different bounds and formatting. The canonical output for two ranges that represent the same set of values, for example the integer ranges \([1, 7]\) and \([1, 8)\), must be identical. It doesn’t matter which representation you choose to be the canonical one, so long as two equivalent values with different formatted are always mapped to the same value with the same formatting. In addition to adjusting the inclusive/exclusive bounds format, a canonicalization function might round off boundary values, in case the desired step size is larger than what the subtype is capable of storing. For instance, a range type over timestamp could be defined to have a step size of an hour, in which case the canonicalization function would need to round off bounds that weren’t a multiple of an hour, or perhaps throw an error instead.

In addition, any range type that is meant to be used with GiST or SP-GiST indexes should define a subtype difference, or `subtype_diff`, function. (The index will still work without `subtype_diff`, but it is likely to be considerably less efficient than if a difference function is provided.) The subtype difference function takes two input values of the subtype, and returns their difference (i.e., \(x\) minus \(y\)) represented as a float8 value. In our example above, the function `float8mi` that underlies the regular float8 minus operator can be used; but for any other subtype, some type conversion would be necessary. Some creative thought about how to represent differences as numbers might be needed, too. To the greatest extent possible, the `subtype_diff` function should agree with the sort ordering implied by the selected operator class and collation; that is, its result should be positive whenever its first argument is greater than its second according to the sort ordering.
A less-oversimplified example of a `subtype_diff` function is:

```
CREATE FUNCTION time_subtype_diff(x time, y time) RETURNS float8 AS
  'SELECT EXTRACT(EPOCH FROM (x - y))'
  LANGUAGE sql STRICT IMMUTABLE;
```

```
CREATE TYPE timerange AS RANGE (
  subtype = time,
  subtype_diff = time_subtype_diff
);
```

```
SELECT '{[11:10, 23:00]}'::timerange;
```

See CREATE TYPE for more information about creating range types.

### 8.17.9. Indexing

GiST and SP-GiST indexes can be created for table columns of range types. For instance, to create a GiST index:

```
CREATE INDEX reservation_idx ON reservation USING GIST (during);
```

A GiST or SP-GiST index can accelerate queries involving these range operators: `=`, `&&`, `<@`, `@>`, `@>>`, `@<<`, `-|`, `&<`, and `&>` (see Table 9-49 for more information).

In addition, B-tree and hash indexes can be created for table columns of range types. For these index types, basically the only useful range operation is equality. There is a B-tree sort ordering defined for range values, with corresponding `<` and `>` operators, but the ordering is rather arbitrary and not usually useful in the real world. Range types’ B-tree and hash support is primarily meant to allow sorting and hashing internally in queries, rather than creation of actual indexes.

### 8.17.10. Constraints on Ranges

While `UNIQUE` is a natural constraint for scalar values, it is usually unsuitable for range types. Instead, an exclusion constraint is often more appropriate (see CREATE TABLE ... CONSTRAINT ... EXCLUDE). Exclusion constraints allow the specification of constraints such as “non-overlapping” on a range type. For example:

```
CREATE TABLE reservation (
  during tsrange,
  EXCLUDE USING GIST (during WITH &&)
);
```

That constraint will prevent any overlapping values from existing in the table at the same time:

```
INSERT INTO reservation VALUES
  ('[2010-01-01 11:30, 2010-01-01 15:00]');
```

```
INSERT INTO reservation VALUES
  ('[2010-01-01 14:45, 2010-01-01 15:45]');
```

```
ERROR: conflicting key value violates exclusion constraint "reservation_during_excl"
DETAIL: Key (during)=('2010-01-01 14:45:00,"2010-01-01 15:45:00") conflicts
with existing key (during)=('2010-01-01 11:30:00,"2010-01-01 15:00:00")
```

180
You can use the `btree_gist` extension to define exclusion constraints on plain scalar data types, which can then be combined with range exclusions for maximum flexibility. For example, after `btree_gist` is installed, the following constraint will reject overlapping ranges only if the meeting room numbers are equal:

```sql
CREATE EXTENSION btree_gist;
CREATE TABLE room_reservation (  
  room text,  
  during tsrange,  
  EXCLUDE USING GIST (room WITH =, during WITH &&) 
);

INSERT INTO room_reservation VALUES
  ('123A', '2010-01-01 14:00, 2010-01-01 15:00');
INSERT 0 1

INSERT INTO room_reservation VALUES
  ('123A', '2010-01-01 14:30, 2010-01-01 15:30');
ERROR: conflicting key value violates exclusion constraint "room_reservation_room_during_excl"
DETAIL: Key (room, during)=(123A, "2010-01-01 14:30:00","2010-01-01 15:30:00") conflicts
with existing key (room, during)=(123A, "2010-01-01 14:00:00","2010-01-01 15:00:00").

INSERT INTO room_reservation VALUES
  ('123B', '2010-01-01 14:30, 2010-01-01 15:30');
INSERT 0 1
```

### 8.18. Object Identifier Types

Object identifiers (OIDs) are used internally by PostgreSQL as primary keys for various system tables. OIDs are not added to user-created tables, unless `WITH OIDS` is specified when the table is created, or the default_with_oids configuration variable is enabled. Type `oid` represents an object identifier. There are also several alias types for `oid`: `regproc`, `regprocedure`, `regoper`, `regoperator`, `regclass`, `regtype`, `regrole`, `regnamespace`, `regconfig`, and `regdictionary`. Table 8-24 shows an overview.

The `oid` type is currently implemented as an unsigned four-byte integer. Therefore, it is not large enough to provide database-wide uniqueness in large databases, or even in large individual tables. So, using a user-created table’s OID column as a primary key is discouraged. OIDs are best used only for references to system tables.

The `oid` type itself has few operations beyond comparison. It can be cast to integer, however, and then manipulated using the standard integer operators. (Beware of possible signed-versus-unsigned confusion if you do this.)

The OID alias types have no operations of their own except for specialized input and output routines. These routines are able to accept and display symbolic names for system objects, rather than the raw numeric value that type `oid` would use. The alias types allow simplified lookup of OID values for objects. For example, to examine the `pg_attribute` rows related to a table `mytable`, one could write:

```sql
SELECT * FROM pg_attribute WHERE attrelid = 'mytable':regclass;
```
rather than:

```
SELECT * FROM pg_attribute
WHERE attrelid = (SELECT oid FROM pg_class WHERE relname = 'mytable');
```

While that doesn’t look all that bad by itself, it’s still oversimplified. A far more complicated sub-select would be needed to select the right OID if there are multiple tables named mytable in different schemas. The regclass input converter handles the table lookup according to the schema path setting, and so it does the “right thing” automatically. Similarly, casting a table’s OID to regclass is handy for symbolic display of a numeric OID.

### Table 8-24. Object Identifier Types

<table>
<thead>
<tr>
<th>Name</th>
<th>References</th>
<th>Description</th>
<th>Value Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>any</td>
<td>numeric object identifier</td>
<td>564182</td>
</tr>
<tr>
<td>regproc</td>
<td>pg_proc</td>
<td>function name</td>
<td>sum</td>
</tr>
<tr>
<td>regprocedure</td>
<td>pg_proc</td>
<td>function with argument types</td>
<td>sum(int4)</td>
</tr>
<tr>
<td>regoperator</td>
<td>pg_operator</td>
<td>operator name</td>
<td>+</td>
</tr>
<tr>
<td>regoperator</td>
<td>pg_operator</td>
<td>operator with argument types</td>
<td>*(integer,integer) or -(NONE,integer)</td>
</tr>
<tr>
<td>regclass</td>
<td>pg_class</td>
<td>relation name</td>
<td>pg_type</td>
</tr>
<tr>
<td>regtype</td>
<td>pg_type</td>
<td>data type name</td>
<td>integer</td>
</tr>
<tr>
<td>regrole</td>
<td>pg_authid</td>
<td>role name</td>
<td>smithee</td>
</tr>
<tr>
<td>regnamespace</td>
<td>pg_namespace</td>
<td>namespace name</td>
<td>pg_catalog</td>
</tr>
<tr>
<td>regconfig</td>
<td>pg_ts_config</td>
<td>text search configuration</td>
<td>english</td>
</tr>
<tr>
<td>regdictionary</td>
<td>pg_ts_dict</td>
<td>text search dictionary</td>
<td>simple</td>
</tr>
</tbody>
</table>

All of the OID alias types for objects grouped by namespace accept schema-qualified names, and will display schema-qualified names on output if the object would not be found in the current search path without being qualified. The regproc and regoperator alias types will only accept input names that are unique (not overloaded), so they are of limited use; for most uses regprocedure or regoperator are more appropriate. For regoperator, unary operators are identified by writing NONE for the unused operand.

An additional property of most of the OID alias types is the creation of dependencies. If a constant of one of these types appears in a stored expression (such as a column default expression or view), it creates a dependency on the referenced object. For example, if a column has a default expression `nextval('my_seq'::regclass)`, PostgreSQL understands that the default expression depends on the sequence my_seq; the system will not let the sequence be dropped without first removing the default expression. regrole is the only exception for the property. Constants of this type are not allowed in such expressions.

**Note:** The OID alias types do not completely follow transaction isolation rules. The planner also treats them as simple constants, which may result in sub-optimal planning.

Another identifier type used by the system is xid, or transaction (abbreviated xact) identifier. This is the data type of the system columns xmin and xmax. Transaction identifiers are 32-bit quantities.
A third identifier type used by the system is cid, or command identifier. This is the data type of the system columns cmin and cmax. Command identifiers are also 32-bit quantities.

A final identifier type used by the system is tid, or tuple identifier (row identifier). This is the data type of the system column ctid. A tuple ID is a pair (block number, tuple index within block) that identifies the physical location of the row within its table.

(The system columns are further explained in Section 5.4.)

8.19. pg_lsn Type

The pg_lsn data type can be used to store LSN (Log Sequence Number) data which is a pointer to a location in the XLOG. This type is a representation of XLogRecPtr and an internal system type of PostgreSQL.

Internally, an LSN is a 64-bit integer, representing a byte position in the write-ahead log stream. It is printed as two hexadecimal numbers of up to 8 digits each, separated by a slash; for example, 16/B374D848. The pg_lsn type supports the standard comparison operators, like = and >. Two LSNs can be subtracted using the - operator; the result is the number of bytes separating those write-ahead log positions.

8.20. Pseudo-Types

The PostgreSQL type system contains a number of special-purpose entries that are collectively called pseudo-types. A pseudo-type cannot be used as a column data type, but it can be used to declare a function’s argument or result type. Each of the available pseudo-types is useful in situations where a function’s behavior does not correspond to simply taking or returning a value of a specific SQL data type. Table 8-25 lists the existing pseudo-types.

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>any</td>
<td>Indicates that a function accepts any input data type.</td>
</tr>
<tr>
<td>anyelement</td>
<td>Indicates that a function accepts any data type (see Section 36.2.5).</td>
</tr>
<tr>
<td>anyarray</td>
<td>Indicates that a function accepts any array data type (see Section 36.2.5).</td>
</tr>
<tr>
<td>anynonarray</td>
<td>Indicates that a function accepts any non-array data type (see Section 36.2.5).</td>
</tr>
<tr>
<td>anyenum</td>
<td>Indicates that a function accepts any enum data type (see Section 36.2.5 and Section 8.7).</td>
</tr>
<tr>
<td>anyrange</td>
<td>Indicates that a function accepts any range data type (see Section 36.2.5 and Section 8.17).</td>
</tr>
<tr>
<td>cstring</td>
<td>Indicates that a function accepts or returns a null-terminated C string.</td>
</tr>
<tr>
<td>internal</td>
<td>Indicates that a function accepts or returns a server-internal data type.</td>
</tr>
</tbody>
</table>
Chapter 8. Data Types

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>language_handler</td>
<td>A procedural language call handler is declared to return language_handler.</td>
</tr>
<tr>
<td>fdw_handler</td>
<td>A foreign-data wrapper handler is declared to return fdw_handler.</td>
</tr>
<tr>
<td>index_am_handler</td>
<td>An index access method handler is declared to return index_am_handler.</td>
</tr>
<tr>
<td>tsm_handler</td>
<td>A tablesample method handler is declared to return tsm_handler.</td>
</tr>
<tr>
<td>record</td>
<td>Identifies a function taking or returning an unspecified row type.</td>
</tr>
<tr>
<td>trigger</td>
<td>A trigger function is declared to return trigger.</td>
</tr>
<tr>
<td>event_trigger</td>
<td>An event trigger function is declared to return event_trigger.</td>
</tr>
<tr>
<td>pg_ddl_command</td>
<td>Identifies a representation of DDL commands that is available to event triggers.</td>
</tr>
<tr>
<td>void</td>
<td>Indicates that a function returns no value.</td>
</tr>
<tr>
<td>opaque</td>
<td>An obsolete type name that formerly served all the above purposes.</td>
</tr>
</tbody>
</table>

Functions coded in C (whether built-in or dynamically loaded) can be declared to accept or return any of these pseudo data types. It is up to the function author to ensure that the function will behave safely when a pseudo-type is used as an argument type.

Functions coded in procedural languages can use pseudo-types only as allowed by their implementation languages. At present most procedural languages forbid use of a pseudo-type as an argument type, and allow only void and record as a result type (plus trigger or event_trigger when the function is used as a trigger or event trigger). Some also support polymorphic functions using the types anyelement, anyarray, anynonarray, anyenum, and anyrange.

The internal pseudo-type is used to declare functions that are meant only to be called internally by the database system, and not by direct invocation in an SQL query. If a function has at least one internal-type argument then it cannot be called from SQL. To preserve the type safety of this restriction it is important to follow this coding rule: do not create any function that is declared to return internal unless it has at least one internal argument.
Chapter 9. Functions and Operators

PostgreSQL provides a large number of functions and operators for the built-in data types. Users can also define their own functions and operators, as described in Part V. The psql commands `\df` and `\do` can be used to list all available functions and operators, respectively.

If you are concerned about portability then note that most of the functions and operators described in this chapter, with the exception of the most trivial arithmetic and comparison operators and some explicitly marked functions, are not specified by the SQL standard. Some of this extended functionality is present in other SQL database management systems, and in many cases this functionality is compatible and consistent between the various implementations. This chapter is also not exhaustive; additional functions appear in relevant sections of the manual.

9.1. Logical Operators

The usual logical operators are available:

- `AND`  
- `OR`  
- `NOT`

SQL uses a three-valued logic system with true, false, and `null`, which represents “unknown”. Observe the following truth tables:

<table>
<thead>
<tr>
<th>a</th>
<th>b</th>
<th>a AND b</th>
<th>a OR b</th>
</tr>
</thead>
<tbody>
<tr>
<td>TRUE</td>
<td>TRUE</td>
<td>TRUE</td>
<td>TRUE</td>
</tr>
<tr>
<td>TRUE</td>
<td>FALSE</td>
<td>FALSE</td>
<td>TRUE</td>
</tr>
<tr>
<td>TRUE</td>
<td>NULL</td>
<td>NULL</td>
<td>TRUE</td>
</tr>
<tr>
<td>FALSE</td>
<td>FALSE</td>
<td>FALSE</td>
<td>FALSE</td>
</tr>
<tr>
<td>FALSE</td>
<td>NULL</td>
<td>FALSE</td>
<td>NULL</td>
</tr>
<tr>
<td>NULL</td>
<td>NULL</td>
<td>NULL</td>
<td>NULL</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>a</th>
<th>NOT a</th>
</tr>
</thead>
<tbody>
<tr>
<td>TRUE</td>
<td>FALSE</td>
</tr>
<tr>
<td>FALSE</td>
<td>TRUE</td>
</tr>
<tr>
<td>NULL</td>
<td>NULL</td>
</tr>
</tbody>
</table>

The operators `AND` and `OR` are commutative, that is, you can switch the left and right operand without affecting the result. But see Section 4.2.14 for more information about the order of evaluation of subexpressions.

9.2. Comparison Functions and Operators

The usual comparison operators are available, as shown in Table 9-1.
Table 9-1. Comparison Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;</td>
<td>less than</td>
</tr>
<tr>
<td>&gt;</td>
<td>greater than</td>
</tr>
<tr>
<td>&lt;=</td>
<td>less than or equal to</td>
</tr>
<tr>
<td>=&gt;</td>
<td>greater than or equal to</td>
</tr>
<tr>
<td>=</td>
<td>equal</td>
</tr>
<tr>
<td>&lt;&gt; or !=</td>
<td>not equal</td>
</tr>
</tbody>
</table>

**Note:** The != operator is converted to <> in the parser stage. It is not possible to implement != and <> operators that do different things.

Comparison operators are available for all relevant data types. All comparison operators are binary operators that return values of type boolean; expressions like 1 < 2 < 3 are not valid (because there is no < operator to compare a Boolean value with 3).

There are also some comparison predicates, as shown in Table 9-2. These behave much like operators, but have special syntax mandated by the SQL standard.

Table 9-2. Comparison Predicates

<table>
<thead>
<tr>
<th>Predicate</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>a BETWEEN x AND y</td>
<td>between</td>
</tr>
<tr>
<td>a NOT BETWEEN x AND y</td>
<td>not between</td>
</tr>
<tr>
<td>a BETWEEN SYMMETRIC x AND y</td>
<td>between, after sorting the comparison values</td>
</tr>
<tr>
<td>a NOT BETWEEN SYMMETRIC x AND y</td>
<td>not between, after sorting the comparison values</td>
</tr>
<tr>
<td>a IS DISTINCT FROM b</td>
<td>not equal, treating null like an ordinary value</td>
</tr>
<tr>
<td>a IS NOT DISTINCT FROM b</td>
<td>equal, treating null like an ordinary value</td>
</tr>
<tr>
<td>expression IS NULL</td>
<td>is null</td>
</tr>
<tr>
<td>expression IS NOT NULL</td>
<td>is not null</td>
</tr>
<tr>
<td>expression ISNULL</td>
<td>is null (nonstandard syntax)</td>
</tr>
<tr>
<td>expression NOTNULL</td>
<td>is not null (nonstandard syntax)</td>
</tr>
<tr>
<td>boolean_expression IS TRUE</td>
<td>is true</td>
</tr>
<tr>
<td>boolean_expression IS NOT TRUE</td>
<td>is false or unknown</td>
</tr>
<tr>
<td>boolean_expression IS FALSE</td>
<td>is false</td>
</tr>
<tr>
<td>boolean_expression IS NOT FALSE</td>
<td>is true or unknown</td>
</tr>
<tr>
<td>boolean_expression IS UNKNOWN</td>
<td>is unknown</td>
</tr>
<tr>
<td>boolean_expression IS NOT UNKNOWN</td>
<td>is true or false</td>
</tr>
</tbody>
</table>

The BETWEEN predicate simplifies range tests:

\[
a \text{ BETWEEN } x \text{ AND } y
\]

is equivalent to

\[
a \geq x \text{ AND } a \leq y
\]
Chapter 9. Functions and Operators

Notice that BETWEEN treats the endpoint values as included in the range. NOT BETWEEN does the opposite comparison:

\[ a \text{ NOT BETWEEN } x \text{ AND } y \]

is equivalent to

\[ a < x \text{ OR } a > y \]

BETWEEN SYMMETRIC is like BETWEEN except there is no requirement that the argument to the left of AND be less than or equal to the argument on the right. If it is not, those two arguments are automatically swapped, so that a nonempty range is always implied.

Ordinary comparison operators yield null (signifying “unknown”), not true or false, when either input is null. For example, \( 7 = \text{NULL} \) yields null, as does \( 7 < > \text{NULL} \). When this behavior is not suitable, use the IS [ NOT ] DISTINCT FROM predicates:

\[ a \text{ IS DISTINCT FROM } b \]
\[ a \text{ IS NOT DISTINCT FROM } b \]

For non-null inputs, IS DISTINCT FROM is the same as the \( <> \) operator. However, if both inputs are null it returns false, and if only one input is null it returns true. Similarly, IS NOT DISTINCT FROM is identical to = for non-null inputs, but it returns true when both inputs are null, and false when only one input is null. Thus, these predicates effectively act as though null were a normal data value, rather than “unknown”.

To check whether a value is or is not null, use the predicates:

\[ \text{expression IS NULL} \]
\[ \text{expression IS NOT NULL} \]

or the equivalent, but nonstandard, predicates:

\[ \text{expression ISNULL} \]
\[ \text{expression NOTNULL} \]

Do not write \( \text{expression} = \text{NULL} \) because NULL is not “equal to” NULL. (The null value represents an unknown value, and it is not known whether two unknown values are equal.)

**Tip:** Some applications might expect that \( \text{expression} = \text{NULL} \) returns true if \( \text{expression} \) evaluates to the null value. It is highly recommended that these applications be modified to comply with the SQL standard. However, if that cannot be done the transform_null_equals configuration variable is available. If it is enabled, PostgreSQL will convert \( x = \text{NULL} \) clauses to \( x \text{ IS NULL} \).

If the \( \text{expression} \) is row-valued, then IS NULL is true when the row expression itself is null or when all the row’s fields are null, while IS NOT NULL is true when the row expression itself is non-null and all the row’s fields are non-null. Because of this behavior, IS NULL and IS NOT NULL do not always return inverse results for row-valued expressions; in particular, a row-valued expression that contains both null and non-null fields will return false for both tests. In some cases, it may be preferable to write \( \text{row IS DISTINCT FROM NULL} \) or \( \text{row IS NOT DISTINCT FROM NULL} \), which will simply check whether the overall row value is null without any additional tests on the row fields.

Boolean values can also be tested using the predicates

\[ \text{boolean_expression IS TRUE} \]
These will always return true or false, never a null value, even when the operand is null. A null input is treated as the logical value "unknown". Notice that IS UNKNOWN and IS NOT UNKNOWN are effectively the same as IS NULL and IS NOT NULL, respectively, except that the input expression must be of Boolean type.

Some comparison-related functions are also available, as shown in Table 9-3.

Table 9-3. Comparison Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>num_nonnulls(VARIADIC &quot;any&quot;)</td>
<td>returns the number of non-null arguments</td>
<td>num_nonnulls(1, NULL, 2)</td>
<td>2</td>
</tr>
<tr>
<td>num_nulls(VARIADIC &quot;any&quot;)</td>
<td>returns the number of null arguments</td>
<td>num_nulls(1, NULL, 2)</td>
<td>1</td>
</tr>
</tbody>
</table>

9.3. Mathematical Functions and Operators

Mathematical operators are provided for many PostgreSQL types. For types without standard mathematical conventions (e.g., date/time types) we describe the actual behavior in subsequent sections. Table 9-4 shows the available mathematical operators.

Table 9-4. Mathematical Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>addition</td>
<td>2 + 3</td>
<td>5</td>
</tr>
<tr>
<td>-</td>
<td>subtraction</td>
<td>2 - 3</td>
<td>-1</td>
</tr>
<tr>
<td>*</td>
<td>multiplication</td>
<td>2 * 3</td>
<td>6</td>
</tr>
<tr>
<td>/</td>
<td>division (integer division truncates the result)</td>
<td>4 / 2</td>
<td>2</td>
</tr>
<tr>
<td>%</td>
<td>modulo (remainder)</td>
<td>5 % 4</td>
<td>1</td>
</tr>
<tr>
<td>^</td>
<td>exponentiation (associates left to right)</td>
<td>2.0 ^ 3.0</td>
<td>8</td>
</tr>
<tr>
<td>\</td>
<td>square root</td>
<td>\ 25.0</td>
<td>5</td>
</tr>
<tr>
<td></td>
<td></td>
<td>cube root</td>
<td></td>
</tr>
<tr>
<td>!</td>
<td>factorial</td>
<td>5 !</td>
<td>120</td>
</tr>
<tr>
<td>!!</td>
<td>factorial (prefix operator)</td>
<td>!! 5</td>
<td>120</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>@</td>
<td>absolute value</td>
<td>@ -5.0</td>
<td>5</td>
</tr>
<tr>
<td>&amp;</td>
<td>bitwise AND</td>
<td>91 &amp; 15</td>
<td>11</td>
</tr>
<tr>
<td></td>
<td></td>
<td>bitwise OR</td>
<td>32</td>
</tr>
<tr>
<td>#</td>
<td>bitwise XOR</td>
<td>17 # 5</td>
<td>20</td>
</tr>
<tr>
<td>~</td>
<td>bitwise NOT</td>
<td>~1</td>
<td>-2</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>bitwise shift left</td>
<td>1 &lt;&lt; 4</td>
<td>16</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>bitwise shift right</td>
<td>8 &gt;&gt; 2</td>
<td>2</td>
</tr>
</tbody>
</table>

The bitwise operators work only on integral data types, whereas the others are available for all numeric data types. The bitwise operators are also available for the bit string types `bit` and `bit varying`, as shown in Table 9-13.

Table 9-5 shows the available mathematical functions. In the table, `dp` indicates double precision. Many of these functions are provided in multiple forms with different argument types. Except where noted, any given form of a function returns the same data type as its argument. The functions working with double precision data are mostly implemented on top of the host system’s C library; accuracy and behavior in boundary cases can therefore vary depending on the host system.

### Table 9-5. Mathematical Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>abs(x)</code></td>
<td>(same as input)</td>
<td>absolute value</td>
<td><code>abs(-17.4)</code></td>
<td>17.4</td>
</tr>
<tr>
<td><code>cbrt(dp)</code></td>
<td>dp</td>
<td>cube root</td>
<td><code>cbrt(27.0)</code></td>
<td>3</td>
</tr>
<tr>
<td><code>ceil(dp or numeric)</code></td>
<td>(same as input)</td>
<td>nearest integer greater than or equal to argument</td>
<td><code>ceil(-42.8)</code></td>
<td>-42</td>
</tr>
<tr>
<td><code>ceiling(dp or numeric)</code></td>
<td>(same as input)</td>
<td>nearest integer greater than or equal to argument</td>
<td><code>ceiling(-95.3)</code></td>
<td>-95</td>
</tr>
<tr>
<td><code>degrees(dp)</code></td>
<td>dp</td>
<td>radians to degrees</td>
<td><code>degrees(0.5)</code></td>
<td>28.6478897565412</td>
</tr>
<tr>
<td><code>div(y numeric, x numeric)</code></td>
<td>numeric</td>
<td>integer quotient of y/x</td>
<td><code>div(9,4)</code></td>
<td>2</td>
</tr>
<tr>
<td><code>exp(dp or numeric)</code></td>
<td>(same as input)</td>
<td>exponential</td>
<td><code>exp(1.0)</code></td>
<td>2.71828182845905</td>
</tr>
<tr>
<td><code>floor(dp or numeric)</code></td>
<td>(same as input)</td>
<td>nearest integer less than or equal to argument</td>
<td><code>floor(-42.8)</code></td>
<td>-43</td>
</tr>
<tr>
<td><code>ln(dp or numeric)</code></td>
<td>(same as input)</td>
<td>natural logarithm</td>
<td><code>ln(2.0)</code></td>
<td>0.693147180559945</td>
</tr>
<tr>
<td><code>log(dp or numeric)</code></td>
<td>(same as input)</td>
<td>base 10 logarithm</td>
<td><code>log(100.0)</code></td>
<td>2</td>
</tr>
<tr>
<td><code>log(b numeric, x numeric)</code></td>
<td>numeric</td>
<td>logarithm to base b</td>
<td><code>log(2, 64.0)</code></td>
<td>6.0000000000</td>
</tr>
<tr>
<td><code>mod(y, x)</code></td>
<td>(same as argument types)</td>
<td>remainder of y/x</td>
<td><code>mod(9,4)</code></td>
<td>1</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>-------------</td>
<td>---------</td>
<td>--------</td>
</tr>
<tr>
<td>pi()</td>
<td>dp</td>
<td>“π” constant</td>
<td>pi()</td>
<td>3.14159265358979</td>
</tr>
<tr>
<td>power(a dp, b dp)</td>
<td>dp</td>
<td>a raised to the power of b</td>
<td>power(9.0, 3.0)</td>
<td>729</td>
</tr>
<tr>
<td>power(a numeric, b numeric)</td>
<td>numeric</td>
<td>a raised to the power of b</td>
<td>power(9.0, 3.0)</td>
<td>729</td>
</tr>
<tr>
<td>radians(dp)</td>
<td>dp</td>
<td>degrees to radians</td>
<td>radians(45.0)</td>
<td>0.785398163397448</td>
</tr>
<tr>
<td>round(dp or numeric)</td>
<td>(same as input)</td>
<td>round to nearest integer</td>
<td>round(42.4)</td>
<td>42</td>
</tr>
<tr>
<td>round(v numeric, s int)</td>
<td>numeric</td>
<td>round to s decimal places</td>
<td>round(42.4382, 2)</td>
<td>42.44</td>
</tr>
<tr>
<td>scale(numeric)</td>
<td>integer</td>
<td>scale of the argument (the number of decimal digits in the fractional part)</td>
<td>scale(8.41)</td>
<td>2</td>
</tr>
<tr>
<td>sign(dp or numeric)</td>
<td>(same as input)</td>
<td>sign of the argument (-1, 0, +1)</td>
<td>sign(-8.4)</td>
<td>-1</td>
</tr>
<tr>
<td>sqrt(dp or numeric)</td>
<td>(same as input)</td>
<td>square root</td>
<td>sqrt(2.0)</td>
<td>1.4142135623731</td>
</tr>
<tr>
<td>trunc(dp or numeric)</td>
<td>(same as input)</td>
<td>truncate toward zero</td>
<td>trunc(42.8)</td>
<td>42</td>
</tr>
<tr>
<td>trunc(v numeric, s int)</td>
<td>numeric</td>
<td>truncate to s decimal places</td>
<td>trunc(42.4382, 2)</td>
<td>42.43</td>
</tr>
<tr>
<td>width_bucket(operand dp, b1 dp, b2 dp, count int)</td>
<td>int</td>
<td>return the bucket number to which operand would be assigned in a histogram having count equal-width buckets spanning the range b1 to b2; returns 0 or count+1 for an input outside the range</td>
<td>width_bucket(5.35, 0.024, 10.06, 5)</td>
<td>35</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

Function | Return Type | Description | Example | Result
---|---|---|---|---
width_bucket(operand numeric, b1 numeric, b2 numeric, count int) | int | return the bucket number to which operand would be assigned in a histogram having count equal-width buckets spanning the range b1 to b2; returns 0 or count+1 for an input outside the range | width_bucket(5.0, 0.024, 10.06, 5) | 3

Table 9-6 shows functions for generating random numbers.

Table 9-6. Random Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>random()</td>
<td>dp</td>
<td>random value in the range $0.0 \leq x &lt; 1.0$</td>
</tr>
<tr>
<td>setseed(dp)</td>
<td>void</td>
<td>set seed for subsequent random() calls (value between -1.0 and 1.0, inclusive)</td>
</tr>
</tbody>
</table>

The characteristics of the values returned by random() depend on the system implementation. It is not suitable for cryptographic applications; see pgcrypto module for an alternative.

Finally, Table 9-7 shows the available trigonometric functions. All trigonometric functions take arguments and return values of type double precision. Each of the trigonometric functions comes in two variants, one that measures angles in radians and one that measures angles in degrees.

Table 9-7. Trigonometric Functions
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function (radians)</th>
<th>Function (degrees)</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>( \text{acos} (x) )</td>
<td>( \text{acosd} (x) )</td>
<td>inverse cosine</td>
</tr>
<tr>
<td>( \text{asin} (x) )</td>
<td>( \text{asind} (x) )</td>
<td>inverse sine</td>
</tr>
<tr>
<td>( \text{atan} (x) )</td>
<td>( \text{atand} (x) )</td>
<td>inverse tangent</td>
</tr>
<tr>
<td>( \text{atan2} (y, x) )</td>
<td>( \text{atan2d} (y, x) )</td>
<td>inverse tangent of ( y/x )</td>
</tr>
<tr>
<td>( \text{cos} (x) )</td>
<td>( \text{cosd} (x) )</td>
<td>cosine</td>
</tr>
<tr>
<td>( \text{cot} (x) )</td>
<td>( \text{cotd} (x) )</td>
<td>cotangent</td>
</tr>
<tr>
<td>( \text{sin} (x) )</td>
<td>( \text{sind} (x) )</td>
<td>sine</td>
</tr>
<tr>
<td>( \text{tan} (x) )</td>
<td>( \text{tand} (x) )</td>
<td>tangent</td>
</tr>
</tbody>
</table>

**Note:** Another way to work with angles measured in degrees is to use the unit transformation functions \( \text{radians()} \) and \( \text{degrees()} \) shown earlier. However, using the degree-based trigonometric functions is preferred, as that way avoids roundoff error for special cases such as \( \text{sind}(30) \).

### 9.4. String Functions and Operators

This section describes functions and operators for examining and manipulating string values. Strings in this context include values of the types character, character varying, and text. Unless otherwise noted, all of the functions listed below work on all of these types, but be wary of potential effects of automatic space-padding when using the character type. Some functions also exist natively for the bit-string types.

SQL defines some string functions that use key words, rather than commas, to separate arguments. Details are in Table 9-8. PostgreSQL also provides versions of these functions that use the regular function invocation syntax (see Table 9-9).

**Note:** Before PostgreSQL 8.3, these functions would silently accept values of several non-string data types as well, due to the presence of implicit coercions from those data types to text. Those coercions have been removed because they frequently caused surprising behaviors. However, the string concatenation operator (||) still accepts non-string input, so long as at least one input is of a string type, as shown in Table 9-8. For other cases, insert an explicit coercion to text if you need to duplicate the previous behavior.

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>string</td>
<td></td>
<td>string</td>
<td>text</td>
<td>String concatenation</td>
</tr>
<tr>
<td>string</td>
<td></td>
<td>non-string or non-string</td>
<td></td>
<td>string</td>
</tr>
<tr>
<td>bit_length(string)</td>
<td>int</td>
<td>Number of bits in string</td>
<td>bit_length('jose')</td>
<td>82</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>-------------------</td>
<td>-------------</td>
<td>---------------------------------</td>
<td>------------------------------</td>
<td>--------</td>
</tr>
<tr>
<td>char_length(string)</td>
<td>int</td>
<td>Number of characters in string</td>
<td>char_length('jose')</td>
<td></td>
</tr>
<tr>
<td>character_length(string)</td>
<td>int</td>
<td>Number of characters in string</td>
<td>character_length('jose')</td>
<td></td>
</tr>
<tr>
<td>lower(string)</td>
<td>text</td>
<td>Convert string to lower case</td>
<td>lower('TOM')</td>
<td>tom</td>
</tr>
<tr>
<td>octet_length(string)</td>
<td>int</td>
<td>Number of bytes in string</td>
<td>octet_length('jose')</td>
<td></td>
</tr>
<tr>
<td>overlay(string placing string from int [for int])</td>
<td>text</td>
<td>Replace substring</td>
<td>overlay('Txxxx' placing 'hom' from 2 for 4)</td>
<td>Thomas</td>
</tr>
<tr>
<td>position(substring in string)</td>
<td>int</td>
<td>Location of specified substring</td>
<td>position('om' in 'Thomas')</td>
<td>3</td>
</tr>
<tr>
<td>substring(string [from int] [for int])</td>
<td>text</td>
<td>Extract substring</td>
<td>substring('Thomas' from 2 for 3)</td>
<td>hom</td>
</tr>
<tr>
<td>substring(string from pattern)</td>
<td>text</td>
<td>Extract substring matching POSIX regular expression. See Section 9.7 for more information on pattern matching.</td>
<td>substring('Thomas' from '...$')</td>
<td>mas</td>
</tr>
<tr>
<td>substring(string from pattern for escape)</td>
<td>text</td>
<td>Extract substring matching SQL regular expression. See Section 9.7 for more information on pattern matching.</td>
<td>substring('Thomas' from '%$o_a#&quot;__' for '#')</td>
<td>mas</td>
</tr>
<tr>
<td>trim([leading</td>
<td>trailing</td>
<td>both] [characters] from string)</td>
<td>text</td>
<td>Remove the longest string containing only characters from characters (a space by default) from the start, end, or both ends (both is the default) of string</td>
</tr>
</tbody>
</table>
Additional string manipulation functions are available and are listed in Table 9-9. Some of them are used internally to implement the SQL-standard string functions listed in Table 9-8.

Table 9-9. Other String Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>trim([leading \</td>
<td>text</td>
<td>Non-standard syntax for trim()</td>
<td>trim(both from 'yxTomxx', 'xyz')</td>
<td>Tom</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>Convert string to upper case</td>
<td>upper('tom')</td>
<td>TOM</td>
</tr>
<tr>
<td></td>
<td>text</td>
<td>int</td>
<td>ASCII code of the first character of the argument. For UTF8 returns the Unicode code point of the character. For other multibyte encodings, the argument must be an ASCII character.</td>
<td>ascii('x')</td>
</tr>
<tr>
<td></td>
<td>text</td>
<td>text</td>
<td>Remove the longest string consisting only of characters in characters (a space by default) from the start and end of string</td>
<td>btrim('xyxtrimyx', 'xyz')</td>
</tr>
</tbody>
</table>

194
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>chr(int)</code></td>
<td>text</td>
<td>Character with the given code. For UTF8 the argument is treated as a Unicode code point. For other multibyte encodings the argument must designate an ASCII character. The NULL (0) character is not allowed because text data types cannot store such bytes.</td>
<td><code>chr(65)</code></td>
<td>A</td>
</tr>
<tr>
<td><code>concat(str &quot;any&quot; [, str &quot;any&quot; [, ...] ] )</code></td>
<td>text</td>
<td>Concatenate the text representations of all the arguments. NULL arguments are ignored.</td>
<td><code>concat('abcde', 2, NULL, 22)</code></td>
<td>abcde222</td>
</tr>
<tr>
<td><code>concat_ws(sep text, str &quot;any&quot; [, str &quot;any&quot; [, ...] ] )</code></td>
<td>text</td>
<td>Concatenate all but the first argument with separators. The first argument is used as the separator string. NULL arguments are ignored.</td>
<td><code>concat_ws(',,', 'abcde', 2, NULL, 22)</code></td>
<td>abcde,2,22</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------------------------</td>
<td>-------------</td>
<td>------------------------------------------------------------------------------</td>
<td>----------------------------------</td>
<td>----------------------------------------------------------------------</td>
</tr>
<tr>
<td>convert(string byte,</td>
<td>bytea</td>
<td>Convert string to dest_encoding. The original encoding is specified by src_encoding. The string must be valid in this encoding. Conversions can be defined by CREATE CONVERSION. Also there are some predefined conversions. See Table 9-10 for available conversions.</td>
<td>convert('text_in_utf8', 'UTF8', 'LATIN1')</td>
<td>text in utf8 represented in Latin-1 encoding (ISO 8859-1)</td>
</tr>
<tr>
<td>src_encoding name,</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>dest_encoding name)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>convert_from(string byte,</td>
<td>text</td>
<td>Convert string to the database encoding. The original encoding is specified by src_encoding. The string must be valid in this encoding.</td>
<td>convert_from('text_in_utf8', 'UTF8')</td>
<td>text in utf8, represented in the current database encoding</td>
</tr>
<tr>
<td>src_encoding name)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>convert_to(string text,</td>
<td>bytea</td>
<td>Convert string to dest_encoding.</td>
<td>convert_to('some text', 'UTF8')</td>
<td>some text represented in the UTF8 encoding</td>
</tr>
<tr>
<td>dest_encoding name)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>decode(string text, format</td>
<td>bytea</td>
<td>Decode binary data from textual representation in string. Options for format are same as in encode.</td>
<td>decode('MTIzAAE=', 'base64')</td>
<td>196</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>encode(data bytea, format text)</td>
<td>text</td>
<td>Encode binary data into a textual representation. Supported formats are: base64, hex, escape. escape converts zero bytes and high-bit-set bytes to octal sequences (\nnn) and doubles backslashes.</td>
<td>encode('123\000\001', 'base64')</td>
<td>MTIzAAE=</td>
</tr>
<tr>
<td>format(formatstr text [, formatarg &quot;any&quot; [, ... ]])</td>
<td>text</td>
<td>Format arguments according to a format string. This function is similar to the C function sprintf. See Section 9.4.1.</td>
<td>format('Hello %s, %1$s', 'World')</td>
<td>Hello World, World</td>
</tr>
<tr>
<td>initcap(string)</td>
<td>text</td>
<td>Convert the first letter of each word to upper case and the rest to lower case. Words are sequences of alphanumeric characters separated by non-alphanumeric characters.</td>
<td>initcap('hi THOMAS')</td>
<td>Hi Thomas</td>
</tr>
<tr>
<td>left(str text, n int)</td>
<td>text</td>
<td>Return first n characters in the string. When n is negative, return all but last</td>
<td>left('abcde', 2)</td>
<td>ab</td>
</tr>
<tr>
<td>length(string)</td>
<td>int</td>
<td>Number of characters in string</td>
<td>length('jose')</td>
<td>4</td>
</tr>
<tr>
<td>length(string bytea, encoding name)</td>
<td>int</td>
<td>Number of characters in string in the given encoding. The string must be valid in this encoding.</td>
<td>length('jose', 'UTF8')</td>
<td>4</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>---------------</td>
<td>-------------</td>
<td>------------------------------------------------------------------------------</td>
<td>-------------------</td>
<td>------------</td>
</tr>
<tr>
<td>lpad(string</td>
<td>text</td>
<td>Fill up the string to length length by prepending the characters fill (a</td>
<td>lpad('hi', 5,</td>
<td>xyxhi</td>
</tr>
<tr>
<td>text, length</td>
<td></td>
<td>space by default). If the string is already longer than length then it is</td>
<td>'xy')</td>
<td></td>
</tr>
<tr>
<td>int[, fill</td>
<td></td>
<td>truncated (on the right).</td>
<td></td>
<td></td>
</tr>
<tr>
<td>text])</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ltrim(string</td>
<td>text</td>
<td>Remove the longest string containing only characters from characters (a</td>
<td>ltrim('zzzytest',</td>
<td>test</td>
</tr>
<tr>
<td>text [,</td>
<td></td>
<td>space by default) from the start of string</td>
<td>'xyz')</td>
<td></td>
</tr>
<tr>
<td>characters</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>text])</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>md5(string)</td>
<td>text</td>
<td>Calculates the MD5 hash of string, returning the result in hexadecimal</td>
<td>md5('abc')</td>
<td>900150983cd24fb0</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td>d6963f7d28e17f72</td>
</tr>
</tbody>
</table>
# Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>parse_ident(qualified_identifier, text [, strictmode boolean DEFAULT true ] )</code></td>
<td><code>text[]</code></td>
<td>Split <code>qualified_identifier</code> into an array of identifiers, removing any quoting of individual identifiers. By default, extra characters after the last identifier are considered an error; but if the second parameter is <code>false</code>, then such extra characters are ignored. (This behavior is useful for parsing names for objects like functions.) Note that this function does not truncate over-length identifiers. If you want truncation you can cast the result to <code>name[]</code>.</td>
<td><code>parse_ident('&quot;SomeSchema&quot;.someTable')</code></td>
<td><code>{SomeSchema,sometable}</code></td>
</tr>
<tr>
<td><code>pg_client_encoding()</code></td>
<td><code>name</code></td>
<td>Current client encoding name</td>
<td><code>pg_client_encoding()</code></td>
<td><code>SQL_ASCII</code></td>
</tr>
<tr>
<td><code>quote_ident(string text)</code></td>
<td><code>text</code></td>
<td>Return the given string suitably quoted to be used as an identifier in an SQL statement string. Quotes are added only if necessary (i.e., if the string contains non-identifier characters or would be case-folded). Embedded quotes are properly doubled. See also Example 41-1.</td>
<td><code>quote_ident('Foo bar')</code></td>
<td><code>&quot;Foo bar&quot;</code></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>-------------</td>
<td>---------</td>
<td>--------</td>
</tr>
<tr>
<td><code>quote_literal(string text)</code></td>
<td>text</td>
<td>Return the given string suitably quoted to be used as a string literal in an SQL statement string. Embedded single-quotes and backslashes are properly doubled. Note that <code>quote_literal</code> returns null on null input; if the argument might be null, <code>quote_nullable</code> is often more suitable. See also Example 41-1.</td>
<td><code>quote_literal(‘O’Reilly’)</code></td>
<td>‘O”Reilly’</td>
</tr>
<tr>
<td><code>quote_literal(anyelement)</code></td>
<td>text</td>
<td>Coerce the given value to text and then quote it as a literal. Embedded single-quotes and backslashes are properly doubled.</td>
<td><code>quote_literal(42.5)</code></td>
<td>‘42.5’</td>
</tr>
<tr>
<td><code>quote_nullable(string text)</code></td>
<td>text</td>
<td>Return the given string suitably quoted to be used as a string literal in an SQL statement string; or, if the argument is null, return NULL. Embedded single-quotes and backslashes are properly doubled. See also Example 41-1.</td>
<td><code>quote_nullable(NULL)</code></td>
<td>NULL</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>--------------------------</td>
<td>-------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
<td>-----------------------------</td>
<td>----------------</td>
</tr>
<tr>
<td><code>quote_nullable</code></td>
<td><code>text</code></td>
<td>Coerce the given value to text and then quote it as a literal; or, if the argument is null, return NULL. Embedded single-quotes and backslashes are properly doubled.</td>
<td><code>quote_nullable(42.5)</code></td>
<td>'42.5'</td>
</tr>
<tr>
<td><code>regexp_matches</code></td>
<td><code>setof text[]</code></td>
<td>Return all captured substrings resulting from matching a POSIX regular expression against the string. See Section 9.7.3 for more information.</td>
<td><code>regexp_matches('foobarbequebaz', '(bar)(beque)')</code></td>
<td>{bar,beque}</td>
</tr>
<tr>
<td><code>regexp_replace</code></td>
<td><code>text</code></td>
<td>Replace substring(s) matching a POSIX regular expression. See Section 9.7.3 for more information.</td>
<td><code>regexp_replace('Thomas', '.[mN]a.', 'M')</code></td>
<td>&quot;ThM&quot;</td>
</tr>
<tr>
<td><code>regexp_split_to_array</code></td>
<td><code>text[]</code></td>
<td>Split string using a POSIX regular expression as the delimiter. See Section 9.7.3 for more information.</td>
<td><code>regexp_split_to_array('hello world', '\s+')</code></td>
<td>{hello,world}</td>
</tr>
<tr>
<td><code>regexp_split_to_table</code></td>
<td><code>setof text</code></td>
<td>Split string using a POSIX regular expression as the delimiter. See Section 9.7.3 for more information.</td>
<td><code>regexp_split_to_table('hello world', '\s+')</code></td>
<td>helloworld (2 rows)</td>
</tr>
<tr>
<td><code>repeat</code></td>
<td><code>text</code></td>
<td>Repeat string the specified number of times</td>
<td><code>repeat('Pg', 4)</code></td>
<td>PgPgPgPg</td>
</tr>
</tbody>
</table>

**Chapter 9. Functions and Operators**
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>replace(string text, from text, to text)</td>
<td>text</td>
<td>Replace all occurrences in string of substring from with substring to</td>
<td>replace('abcdef', 'cd', 'XX')</td>
<td>abXXefXXef</td>
</tr>
<tr>
<td>reverse(str)</td>
<td>text</td>
<td>Return reversed string.</td>
<td>reverse('abcded')</td>
<td>dcba</td>
</tr>
<tr>
<td>right(str text, n int)</td>
<td>text</td>
<td>Return last n characters in the string. When n is negative, return all but first</td>
<td>right('abcde', 2)</td>
<td>de</td>
</tr>
<tr>
<td>rpad(string text, length int [, fill text])</td>
<td>text</td>
<td>Fill up the string to length by appending the characters fill (a space by default). If the string is already longer than length then it is truncated.</td>
<td>rpad('hi', 5, 'xy')</td>
<td>hixyx</td>
</tr>
<tr>
<td>rtrim(string text [, characters text])</td>
<td>text</td>
<td>Remove the longest string containing only characters from characters (a space by default) from the end of string</td>
<td>rtrim('testxxzx', 'xyz')</td>
<td>test</td>
</tr>
<tr>
<td>split_part(string text, delimiter text, field int)</td>
<td>text</td>
<td>Split string on delimiter and return the given field (counting from one)</td>
<td>split_part('abcdef-@-ghi', '@@-', 2)</td>
<td>abcdeff</td>
</tr>
<tr>
<td>strpos(string, substring)</td>
<td>int</td>
<td>Location of specified substring (same as position(substring in string), but note the reversed argument order)</td>
<td>strpos('high', 'ig')</td>
<td>2</td>
</tr>
<tr>
<td>substr(string, from [, count])</td>
<td>text</td>
<td>Extract substring (same as substring(string from from for count))</td>
<td>substr('alphabet', 3, 2)</td>
<td>ph</td>
</tr>
</tbody>
</table>
## Functions and Operators

### Table 9-10. Built-in Conversions

<table>
<thead>
<tr>
<th>Conversion Name</th>
<th>Source Encoding</th>
<th>Destination Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>ascii_to_mic</td>
<td>SQL_ASCII</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>ascii_to_utf8</td>
<td>SQL_ASCII</td>
<td>UTF8</td>
</tr>
<tr>
<td>big5_to_euc_tw</td>
<td>EUC_JP</td>
<td>EUC_TW</td>
</tr>
<tr>
<td>big5_to_mic</td>
<td>EUC_JP</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>big5_to_utf8</td>
<td>EUC_JP</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_cn_to_mic</td>
<td>EUC_CN</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_cn_to_utf8</td>
<td>EUC_CN</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_jp_to_mic</td>
<td>EUC_JP</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_jp_to_sjis</td>
<td>EUC_JP</td>
<td>SJIS</td>
</tr>
</tbody>
</table>

The `concat`, `concat_ws` and `format` functions are variadic, so it is possible to pass the values to be concatenated or formatted as an array marked with the `VARIADIC` keyword (see Section 36.4.5). The array’s elements are treated as if they were separate ordinary arguments to the function. If the variadic array argument is NULL, `concat` and `concat_ws` return NULL, but `format` treats a NULL as a zero-element array.

See also the aggregate function `string_agg` in Section 9.20.
## Functions and Operators

<table>
<thead>
<tr>
<th>Conversion Name</th>
<th>Source Encoding</th>
<th>Destination Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>euc_jp_to_utf8</td>
<td>EUC_JP</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_kr_to_mic</td>
<td>EUC_KR</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_kr_to_utf8</td>
<td>EUC_KR</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_tw_to_big5</td>
<td>EUC_TW</td>
<td>BIG5</td>
</tr>
<tr>
<td>euc_tw_to_mic</td>
<td>EUC_TW</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>euc_tw_to_utf8</td>
<td>EUC_TW</td>
<td>UTF8</td>
</tr>
<tr>
<td>gb18030_to_utf8</td>
<td>GB18030</td>
<td>UTF8</td>
</tr>
<tr>
<td>gbk_to_utf8</td>
<td>GBK</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_10_to_utf8</td>
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<tr>
<td>iso_8859_13_to_utf8</td>
<td>LATIN7</td>
<td>UTF8</td>
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<tr>
<td>iso_8859_14_to_utf8</td>
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<td>iso_8859_15_to_utf8</td>
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<td>UTF8</td>
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<tr>
<td>iso_8859_16_to_utf8</td>
<td>LATIN10</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_1_to_mic</td>
<td>LATIN1</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>iso_8859_1_to_utf8</td>
<td>LATIN1</td>
<td>UTF8</td>
</tr>
<tr>
<td>iso_8859_2_to_mic</td>
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<tr>
<td>iso_8859_2_to_utf8</td>
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<td>iso_8859_2_to_windows_1250</td>
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<tr>
<td>iso_8859_3_to_mic</td>
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<td>iso_8859_3_to_utf8</td>
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<td>iso_8859_4_to_utf8</td>
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<td>iso_8859_5_to_koi8_r</td>
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<tr>
<td>iso_8859_5_to_windows_1251</td>
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<tr>
<td>iso_8859_5_to_windows_866</td>
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<td>WIN866</td>
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<tr>
<td>iso_8859_9_to_utf8</td>
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<tr>
<td>johab_to_utf8</td>
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</tr>
<tr>
<td>koi8_r_to_iso_8859_5</td>
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<tr>
<td>mic_to_ascii</td>
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</tr>
<tr>
<td>Conversion Name</td>
<td>Source Encoding</td>
<td>Destination Encoding</td>
</tr>
<tr>
<td>----------------</td>
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<td>----------------------</td>
</tr>
<tr>
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</tr>
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<td>utf8_to_gbk</td>
<td>UTF8</td>
<td>GBK</td>
</tr>
<tr>
<td>utf8_to_iso_8859_1</td>
<td>UTF8</td>
<td>LATIN1</td>
</tr>
<tr>
<td>utf8_to_iso_8859_10</td>
<td>UTF8</td>
<td>LATIN6</td>
</tr>
<tr>
<td>utf8_to_iso_8859_13</td>
<td>UTF8</td>
<td>LATIN7</td>
</tr>
<tr>
<td>utf8_to_iso_8859_14</td>
<td>UTF8</td>
<td>LATIN8</td>
</tr>
<tr>
<td>utf8_to_iso_8859_15</td>
<td>UTF8</td>
<td>LATIN9</td>
</tr>
<tr>
<td>utf8_to_iso_8859_16</td>
<td>UTF8</td>
<td>LATIN10</td>
</tr>
<tr>
<td>utf8_to_iso_8859_2</td>
<td>UTF8</td>
<td>LATIN2</td>
</tr>
<tr>
<td>utf8_to_iso_8859_3</td>
<td>UTF8</td>
<td>LATIN3</td>
</tr>
<tr>
<td>utf8_to_iso_8859_4</td>
<td>UTF8</td>
<td>LATIN4</td>
</tr>
<tr>
<td>utf8_to_iso_8859_5</td>
<td>UTF8</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>utf8_to_iso_8859_6</td>
<td>UTF8</td>
<td>ISO_8859_6</td>
</tr>
<tr>
<td>utf8_to_iso_8859_7</td>
<td>UTF8</td>
<td>ISO_8859_7</td>
</tr>
<tr>
<td>utf8_to_iso_8859_8</td>
<td>UTF8</td>
<td>ISO_8859_8</td>
</tr>
<tr>
<td>utf8_to_iso_8859_9</td>
<td>UTF8</td>
<td>LATIN5</td>
</tr>
</tbody>
</table>

205
<table>
<thead>
<tr>
<th>Conversion Name</th>
<th>Source Encoding</th>
<th>Destination Encoding</th>
</tr>
</thead>
<tbody>
<tr>
<td>utf8_to_johab</td>
<td>UTF8</td>
<td>JOHAB</td>
</tr>
<tr>
<td>utf8_to_koi8_r</td>
<td>UTF8</td>
<td>KOI8R</td>
</tr>
<tr>
<td>utf8_to_koi8_u</td>
<td>UTF8</td>
<td>KOI8U</td>
</tr>
<tr>
<td>utf8_to_sjis</td>
<td>UTF8</td>
<td>SJIS</td>
</tr>
<tr>
<td>utf8_to_windows_1258</td>
<td>UTF8</td>
<td>WIN1258</td>
</tr>
<tr>
<td>utf8_to_uhc</td>
<td>UTF8</td>
<td>UHC</td>
</tr>
<tr>
<td>utf8_to_windows_1250</td>
<td>UTF8</td>
<td>WIN1250</td>
</tr>
<tr>
<td>utf8_to_windows_1251</td>
<td>UTF8</td>
<td>WIN1251</td>
</tr>
<tr>
<td>utf8_to_windows_1252</td>
<td>UTF8</td>
<td>WIN1252</td>
</tr>
<tr>
<td>utf8_to_windows_1253</td>
<td>UTF8</td>
<td>WIN1253</td>
</tr>
<tr>
<td>utf8_to_windows_1254</td>
<td>UTF8</td>
<td>WIN1254</td>
</tr>
<tr>
<td>utf8_to_windows_1255</td>
<td>UTF8</td>
<td>WIN1255</td>
</tr>
<tr>
<td>utf8_to_windows_1256</td>
<td>UTF8</td>
<td>WIN1256</td>
</tr>
<tr>
<td>utf8_to_windows_1257</td>
<td>UTF8</td>
<td>WIN1257</td>
</tr>
<tr>
<td>utf8_to_windows_866</td>
<td>UTF8</td>
<td>WIN866</td>
</tr>
<tr>
<td>utf8_to_windows_874</td>
<td>UTF8</td>
<td>WIN874</td>
</tr>
<tr>
<td>windows_1250_to_iso_8859</td>
<td>WIN1250</td>
<td>LATIN2</td>
</tr>
<tr>
<td>windows_1250_to_mic</td>
<td>WIN1250</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>windows_1250_to_utf8</td>
<td>WIN1250</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_1251_to_iso_8859</td>
<td>WIN1251</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>windows_1251_to_koi8_r</td>
<td>WIN1251</td>
<td>KOI8R</td>
</tr>
<tr>
<td>windows_1251_to_mic</td>
<td>WIN1251</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>windows_1251_to_utf8</td>
<td>WIN1251</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_1251_to_windows_1251</td>
<td>WIN1251</td>
<td>WIN866</td>
</tr>
<tr>
<td>windows_1252_to_utf8</td>
<td>WIN1252</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_1256_to_utf8</td>
<td>WIN1256</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_866_to_iso_8859</td>
<td>WIN866</td>
<td>ISO_8859_5</td>
</tr>
<tr>
<td>windows_866_to_koi8_r</td>
<td>WIN866</td>
<td>KOI8R</td>
</tr>
<tr>
<td>windows_866_to_mic</td>
<td>WIN866</td>
<td>MULE_INTERNAL</td>
</tr>
<tr>
<td>windows_866_to_utf8</td>
<td>WIN866</td>
<td>UTF8</td>
</tr>
<tr>
<td>windows_866_to_windows_1251</td>
<td>WIN866</td>
<td>WIN</td>
</tr>
<tr>
<td>windows_874_to_utf8</td>
<td>WIN874</td>
<td>UTF8</td>
</tr>
<tr>
<td>euc_jis_2004_to_utf8</td>
<td>EUC_JIS_2004</td>
<td>UTF8</td>
</tr>
<tr>
<td>utf8_to_euc_jis_2004</td>
<td>UTF8</td>
<td>EUC_JIS_2004</td>
</tr>
<tr>
<td>shift_jis_2004_to_utf8</td>
<td>SHIFT_JIS_2004</td>
<td>UTF8</td>
</tr>
<tr>
<td>utf8_to_shift_jis_2004</td>
<td>UTF8</td>
<td>SHIFT_JIS_2004</td>
</tr>
<tr>
<td>euc_jis_2004_to_shift_jis</td>
<td>EUC_JIS_2004</td>
<td>SHIFT_JIS_2004</td>
</tr>
</tbody>
</table>
Conversion Name | Source Encoding | Destination Encoding
---|---|---

Notes:
a. The conversion names follow a standard naming scheme: The official name of the source encoding with all non-alphanumeric characters replaced by underscores, followed by_to_, followed by the similarly processed destination encoding name. Therefore, the names might deviate from the customary encoding names.

9.4.1. format

The function format produces output formatted according to a format string, in a style similar to the C function sprintf.

```plaintext
format{formatstr text [, formatarg "any" [, ... ]])
```

`formatstr` is a format string that specifies how the result should be formatted. Text in the format string is copied directly to the result, except where `format specifiers` are used. Format specifiers act as placeholders in the string, defining how subsequent function arguments should be formatted and inserted into the result. Each `formatarg` argument is converted to text according to the usual output rules for its data type, and then formatted and inserted into the result string according to the format specifier(s).

Format specifiers are introduced by a `%` character and have the form

```
%[position][flags][width]type
```

where the component fields are:

**position** (optional)

A string of the form `n$` where `n` is the index of the argument to print. Index 1 means the first argument after `formatstr`. If the `position` is omitted, the default is to use the next argument in sequence.

**flags** (optional)

Additional options controlling how the format specifier’s output is formatted. Currently the only supported flag is a minus sign (`-`) which will cause the format specifier’s output to be left-justified. This has no effect unless the `width` field is also specified.

**width** (optional)

Specifies the minimum number of characters to use to display the format specifier’s output. The output is padded on the left or right (depending on the `-` flag) with spaces as needed to fill the width. A too-small width does not cause truncation of the output, but is simply ignored. The width may be specified using any of the following: a positive integer; an asterisk (`*`) to use the next function argument as the width; or a string of the form `*n$` to use the `n`th function argument as the width.

If the width comes from a function argument, that argument is consumed before the argument that is used for the format specifier’s value. If the width argument is negative, the result is left aligned (as if the `-` flag had been specified) within a field of length `abs(width)`. 
Chapter 9. Functions and Operators

**type (required)**

The type of format conversion to use to produce the format specifier’s output. The following types are supported:

- **s** formats the argument value as a simple string. A null value is treated as an empty string.
- **I** treats the argument value as an SQL identifier, double-quoting it if necessary. It is an error for the value to be null (equivalent to `quote_ident`).
- **L** quotes the argument value as an SQL literal. A null value is displayed as the string NULL, without quotes (equivalent to `quote_nullable`).

In addition to the format specifiers described above, the special sequence `%%` may be used to output a literal `%` character.

Here are some examples of the basic format conversions:

```
SELECT format('Hello %s', 'World');
Result: Hello World

SELECT format('Testing %s, %s, %s, %s', 'one', 'two', 'three');
Result: Testing one, two, three, %

SELECT format('INSERT INTO %I VALUES(%L)', 'Foo bar', E'O\'Reilly');
Result: INSERT INTO "Foo bar" VALUES('O"Reilly')

SELECT format('INSERT INTO %I VALUES(%L)', 'locations', 'C:\Program Files');
Result: INSERT INTO locations VALUES('C:\Program Files')
```

Here are examples using `width` fields and the `-` flag:

```
SELECT format('|%10s|', 'foo');
Result: | foo|

SELECT format('|%-10s|', 'foo');
Result: |foo |

SELECT format('|%s|', 10, 'foo');
Result: | foo|

SELECT format('|%s|', -10, 'foo');
Result: |foo |

SELECT format('|%-s|', 10, 'foo');
Result: |foo |

SELECT format('|%-s|', -10, 'foo');
Result: |foo |
```

These examples show use of `position` fields:

```
SELECT format('Testing %s, %s, %s', 'one', 'two', 'three');
Result: Testing three, two, one
```
Unlike the standard C function `sprintf`, PostgreSQL’s `format` function allows format specifiers with and without position fields to be mixed in the same format string. A format specifier without a position field always uses the next argument after the last argument consumed. In addition, the `format` function does not require all function arguments to be used in the format string. For example:

SELECT format('Testing %3$s, %2$s, %s', 'one', 'two', 'three');
Result: Testing three, two, three

The `%I` and `%L` format specifiers are particularly useful for safely constructing dynamic SQL statements. See Example 41-1.

9.5. Binary String Functions and Operators

This section describes functions and operators for examining and manipulating values of type `bytea`.

SQL defines some string functions that use key words, rather than commas, to separate arguments. Details are in Table 9-11. PostgreSQL also provides versions of these functions that use the regular function invocation syntax (see Table 9-12).

Note: The sample results shown on this page assume that the server parameter `bytea_output` is set to `escape` (the traditional PostgreSQL format).

### Table 9-11. SQL Binary String Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
</table>
| string || string | bytea               | String concatenation                                         | 'Postgres\000' |||bytea
| octet_length(string) | int        | Number of bytes in binary string                             | octet_length('jo\002se')::bytea | 5                |
| overlay(string placing string from int [for int]) | bytea   | Replace substring                                           | overlay('Th\0021b\0023mas placing '\002\003::bytea from 2 for 3) | \002\0021b\0023mas
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>position(substring in string)</td>
<td>int</td>
<td>Location of specified substring</td>
<td>position(‘\000\001om’::bytea in ‘Th\000omas’::bytea)</td>
<td></td>
</tr>
<tr>
<td>substring(string [from int] [for int])</td>
<td>bytea</td>
<td>Extract substring</td>
<td>substring(‘Th\000\001om\000\001ae’::bytea from 2 for 3)</td>
<td></td>
</tr>
<tr>
<td>trim([both] bytes from string)</td>
<td>bytea</td>
<td>Remove the longest string containing only bytes appearing in bytes from the start and end of string</td>
<td>trim(‘\000\001trim\001’::bytea from ‘\000T\001om\000\001’::bytea)</td>
<td></td>
</tr>
</tbody>
</table>

Additional binary string manipulation functions are available and are listed in Table 9-12. Some of them are used internally to implement the SQL-standard string functions listed in Table 9-11.

Table 9-12. Other Binary String Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>btrim(string bytea, bytes bytea)</td>
<td>bytea</td>
<td>Remove the longest string containing only bytes appearing in bytes from the start and end of string</td>
<td>btrim(‘\000trim\001’::bytea, ‘\000\001’::bytea)</td>
<td></td>
</tr>
<tr>
<td>decode(string text, format text)</td>
<td>bytea</td>
<td>Decode binary data from textual representation in string. Options for format are same as in encode.</td>
<td>decode(‘123\004\000456’ ‘escape’)</td>
<td></td>
</tr>
<tr>
<td>encode(data bytea, format text)</td>
<td>text</td>
<td>Encode binary data into a textual representation. Supported formats are: base64, hex, escape. escape converts zero bytes and high-bit-set bytes to octal sequences (\nnn) and doubles backslashes.</td>
<td>encode(‘123\004\000456a, ‘escape’)</td>
<td></td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

#### 9.6. Bit String Functions and Operators

This section describes functions and operators for examining and manipulating bit strings, that is values of the types `bit` and `bit varying`. Aside from the usual comparison operators, the operators shown in Table 9-13 can be used. Bit string operands of `&`, `|`, and `#` must be of equal length. When bit shifting, the original length of the string is preserved, as shown in the examples.

**Table 9-13. Bit String Operators**

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>`</td>
<td></td>
<td>`</td>
<td>concatenation</td>
</tr>
<tr>
<td><code>&amp;</code></td>
<td>bitwise AND</td>
<td>B'10001' &amp; B'01101'</td>
<td>00001</td>
</tr>
<tr>
<td>`</td>
<td>`</td>
<td>bitwise OR</td>
<td>B'10001'</td>
</tr>
<tr>
<td><code>#</code></td>
<td>bitwise XOR</td>
<td>B'10001' # B'01101'</td>
<td>11100</td>
</tr>
<tr>
<td><code>~</code></td>
<td>bitwise NOT</td>
<td>~ B'10001'</td>
<td>01110</td>
</tr>
<tr>
<td><code>&lt;&lt;</code></td>
<td>bitwise shift left</td>
<td>B'10001' &lt;&lt; 3</td>
<td>010000</td>
</tr>
</tbody>
</table>
### 9. Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>&gt;&gt;</code></td>
<td>bitwise shift right</td>
<td>B'10001' <code>&gt;&gt;</code> 2</td>
<td>00100</td>
</tr>
</tbody>
</table>

The following SQL-standard functions work on bit strings as well as character strings: `length`, `bit_length`, `octet_length`, `position`, `substring`, `overlay`.

The following functions work on bit strings as well as binary strings: `get_bit`, `set_bit`. When working with a bit string, these functions number the first (leftmost) bit of the string as bit 0.

In addition, it is possible to cast integral values to and from type `bit`. Some examples:

- `44::bit(10)`
- `44::bit(3)`
- `cast(-44 as bit(12))`
- `'1110'::bit(4)::integer`

Note that casting to just “bit” means casting to `bit(1)`, and so will deliver only the least significant bit of the integer.

**Note:** Casting an integer to `bit(n)` copies the rightmost `n` bits. Casting an integer to a bit string width wider than the integer itself will sign-extend on the left.

### 9.7. Pattern Matching

There are three separate approaches to pattern matching provided by PostgreSQL: the traditional SQL `LIKE` operator, the more recent `SIMILAR TO` operator (added in SQL:1999), and POSIX-style regular expressions. Aside from the basic “does this string match this pattern?” operators, functions are available to extract or replace matching substrings and to split a string at matching locations.

**Tip:** If you have pattern matching needs that go beyond this, consider writing a user-defined function in Perl or Tcl.

#### 9.7.1. LIKE

- `string LIKE pattern [ESCAPE escape-character]`
- `string NOT LIKE pattern [ESCAPE escape-character]`
The **LIKE** expression returns true if the *string* matches the supplied *pattern*. (As expected, the **NOT LIKE** expression returns false if **LIKE** returns true, and vice versa. An equivalent expression is **NOT (string LIKE pattern)**.)

If *pattern* does not contain percent signs or underscores, then the pattern only represents the string itself; in that case **LIKE** acts like the equals operator. An underscore (_) in *pattern* stands for (matches) any single character; a percent sign (%) matches any sequence of zero or more characters.

Some examples:

```sql
'abc' LIKE 'abc'  true
'abc' LIKE 'a%'  true
'abc' LIKE '_b_'  true
'abc' LIKE 'c'   false
```

**LIKE** pattern matching always covers the entire string. Therefore, if it’s desired to match a sequence anywhere within a string, the pattern must start and end with a percent sign.

To match a literal underscore or percent sign without matching other characters, the respective character in *pattern* must be preceded by the escape character. The default escape character is the backslash but a different one can be selected by using the **ESCAPE** clause. To match the escape character itself, write two escape characters.

**Note:** If you have **standard_conforming_strings** turned off, any backslashes you write in literal string constants will need to be doubled. See Section 4.1.2.1 for more information.

It’s also possible to select no escape character by writing **ESCAPE "**. This effectively disables the escape mechanism, which makes it impossible to turn off the special meaning of underscore and percent signs in the pattern.

The key word **ILIKE** can be used instead of **LIKE** to make the match case-insensitive according to the active locale. This is not in the SQL standard but is a PostgreSQL extension.

The operator `~~` is equivalent to **LIKE**, and `~~*` corresponds to **ILIKE**. There are also `!~~` and `!~~*` operators that represent **NOT LIKE** and **NOT ILIKE**, respectively. All of these operators are PostgreSQL-specific.

### 9.7.2. **SIMILAR TO** Regular Expressions

```sql
string SIMILAR TO pattern [ESCAPE escape-character]
string NOT SIMILAR TO pattern [ESCAPE escape-character]
```

The **SIMILAR TO** operator returns true or false depending on whether its pattern matches the given string. It is similar to **LIKE**, except that it interprets the pattern using the SQL standard’s definition of a regular expression. SQL regular expressions are a curious cross between **LIKE** notation and common regular expression notation.

Like **LIKE**, the **SIMILAR TO** operator succeeds only if its pattern matches the entire string; this is unlike common regular expression behavior where the pattern can match any part of the string. Also like **LIKE**, **SIMILAR TO** uses _ and % as wildcard characters denoting any single character and any string, respectively (these are comparable to . and .* in POSIX regular expressions).
In addition to these facilities borrowed from LIKE, SIMILAR TO supports these pattern-matching metacharacters borrowed from POSIX regular expressions:

- | denotes alternation (either of two alternatives).
- * denotes repetition of the previous item zero or more times.
- + denotes repetition of the previous item one or more times.
- ? denotes repetition of the previous item zero or one time.
- {m} denotes repetition of the previous item exactly m times.
- {m,} denotes repetition of the previous item m or more times.
- {m,n} denotes repetition of the previous item at least m and not more than n times.
- Parentheses () can be used to group items into a single logical item.
- A bracket expression [ . . ] specifies a character class, just as in POSIX regular expressions.

Notice that the period ( . ) is not a metacharacter for SIMILAR TO.

As with LIKE, a backslash disables the special meaning of any of these metacharacters; or a different escape character can be specified with ESCAPE.

Some examples:

'abc' SIMILAR TO 'abc' true
'abc' SIMILAR TO 'a' false
'abc' SIMILAR TO '%(b|d)%' true
'abc' SIMILAR TO '(b|c)%' false

The substring function with three parameters, substring(string from pattern for escape-character), provides extraction of a substring that matches an SQL regular expression pattern. As with SIMILAR TO, the specified pattern must match the entire data string, or else the function fails and returns null. To indicate the part of the pattern that should be returned on success, the pattern must contain two occurrences of the escape character followed by a double quote (" ). The text matching the portion of the pattern between these markers is returned.

Some examples, with '#' delimiting the return string:

substring('foobar' from '%#'o_b#'%' for '#') oob
substring('foobar' from '#o_b#%' for '#') NULL

9.7.3. POSIX Regular Expressions

Table 9-14 lists the available operators for pattern matching using POSIX regular expressions.

Table 9-14. Regular Expression Match Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>~</td>
<td>Matches regular expression, case sensitive</td>
<td>'thomas' ~ '.<em>thomas.</em>'</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>~*</td>
<td>Matches regular expression, case insensitive</td>
<td>'thomas’ ~* ‘.<em>Thomas.</em>’</td>
</tr>
<tr>
<td>!~</td>
<td>Does not match regular expression, case sensitive</td>
<td>'thomas’ !~ ‘.<em>Thomas.</em>’</td>
</tr>
<tr>
<td>!~*</td>
<td>Does not match regular expression, case insensitive</td>
<td>'thomas’ !~* ‘.<em>vadim.</em>’</td>
</tr>
</tbody>
</table>

POSIX regular expressions provide a more powerful means for pattern matching than the LIKE and SIMILAR TO operators. Many Unix tools such as `egrep`, `sed`, or `awk` use a pattern matching language that is similar to the one described here.

A regular expression is a character sequence that is an abbreviated definition of a set of strings (a regular set). A string is said to match a regular expression if it is a member of the regular set described by the regular expression. As with LIKE, pattern characters match string characters exactly unless they are special characters in the regular expression language — but regular expressions use different special characters than LIKE does. Unlike LIKE patterns, a regular expression is allowed to match anywhere within a string, unless the regular expression is explicitly anchored to the beginning or end of the string.

Some examples:

- `'abc’ ~ ‘abc’` true
- `'abc’ ~ ‘^a’` true
- `'abc’ ~ ‘(b|d)’` true
- `'abc’ ~ ‘^c’` false

The POSIX pattern language is described in much greater detail below.

The `substring` function with two parameters, `substring(string from pattern)`, provides extraction of a substring that matches a POSIX regular expression pattern. It returns null if there is no match, otherwise the portion of the text that matched the pattern. But if the pattern contains any parentheses, the portion of the text that matched the first parenthesized subexpression (the one whose left parenthesis comes first) is returned. You can put parentheses around the whole expression if you want to use parentheses without triggering this exception. If you need parentheses in the pattern before the subexpression you want to extract, see the non-capturing parentheses described below.

Some examples:

- `substring('foobar' from 'o.b')` oob
- `substring('foobar' from 'o(.)b')` o

The `regexp_replace` function provides substitution of new text for substrings that match POSIX regular expression patterns. It has the syntax `regexp_replace(source, pattern, replacement [, flags])`. The source string is returned unchanged if there is no match to the pattern. If there is a match, the source string is returned with the replacement string substituted for the matching substring. The replacement string can contain `\n`, where `n` is 1 through 9, to indicate that the source substring matching the `n`th parenthesized subexpression of the pattern should be inserted, and it can contain `\` to indicate that the substring matching the entire pattern should be inserted. Write `\` if you need to put a literal backslash in the replacement text. The flags parameter is an optional text string containing zero or more single-letter flags that change the function’s behavior. Flag `i` specifies
case-insensitive matching, while flag g specifies replacement of each matching substring rather than
only the first one. Supported flags (though not g) are described in Table 9-22.

Some examples:

```
regexp_replace('foobarbaz', 'b..', 'X')
fooXbaz
regexp_replace('foobarbaz', 'b..', 'X', 'g')
fooXX
regexp_replace('foobarbaz', 'b(..)', 'X\1Y', 'g')
fooXarYXazY
```

The `regexp_matches` function returns a text array of all of the captured substrings resulting from
matching a POSIX regular expression pattern. It has the syntax `regexp_matches(string, pattern [, flags])`. The function can return no rows, one row, or multiple rows (see the g flag below). If the
`pattern` does not match, the function returns no rows. If the pattern contains no parenthesized substrings,
then each row returned is a single-element text array containing the substring matching the whole pattern. If the pattern contains parenthesized subexpressions, the function returns a text array
whose n'th element is the substring matching the n'th parenthesized subexpression of the pattern (not
counting “non-capturing” parentheses; see below for details). The `flags` parameter is an optional text
string containing zero or more single-letter flags that change the function’s behavior. Flag g causes
the function to find each match in the string, not only the first one, and return a row for each such
match. Supported flags (though not g) are described in Table 9-22.

Some examples:

```
SELECT regexp_matches('foobarbequebaz', '(bar)(beque)');
regexp_matches
----------------
{bar,beque}
(1 row)
SELECT regexp_matches('foobaroxyzilbarfbonk', '(b[^b]+)(b[^b]+)', 'g');
regexp_matches
----------------
{bar,beque}
{bazil,barf}
(2 rows)
SELECT regexp_matches('foobaroxyzb', 'barbeque');
regexp_matches
----------------
{barbeque}
(1 row)
```

It is possible to force `regexp_matches()` to always return one row by using a sub-select; this is
particularly useful in a `SELECT` target list when you want all rows returned, even non-matching ones:

```
SELECT col1, (SELECT regexp_matches(col2, '(bar)(beque)')) FROM tab;
```

The `regexp_split_to_table` function splits a string using a POSIX regular expression pattern as
a delimiter. It has the syntax `regexp_split_to_table(string, pattern [, flags])`. If there is
no match to the `pattern`, the function returns the `string`. If there is at least one match, for each
match it returns the text from the end of the last match (or the beginning of the string) to the beginning of the match. When there are no more matches, it returns the text from the end of the last match to the end of the string. The *flags* parameter is an optional text string containing zero or more single-letter flags that change the function’s behavior. *regexp_split_to_table* supports the flags described in Table 9-22.

The *regexp_split_to_array* function behaves the same as *regexp_split_to_table*, except that *regexp_split_to_array* returns its result as an array of text. It has the syntax

```
regexp_split_to_array(string, pattern [, flags ])
```

The parameters are the same as for *regexp_split_to_table*.

Some examples:

```
SELECT foo FROM regexp_split_to_table('the quick brown fox jumps over the lazy dog', '\s+') AS foo;
```

<table>
<thead>
<tr>
<th>foo</th>
</tr>
</thead>
<tbody>
<tr>
<td>the</td>
</tr>
<tr>
<td>quick</td>
</tr>
<tr>
<td>brown</td>
</tr>
<tr>
<td>fox</td>
</tr>
<tr>
<td>jumps</td>
</tr>
<tr>
<td>over</td>
</tr>
<tr>
<td>the</td>
</tr>
<tr>
<td>lazy</td>
</tr>
<tr>
<td>dog</td>
</tr>
</tbody>
</table>

(9 rows)

```
SELECT regexp_split_to_array('the quick brown fox jumps over the lazy dog', '\s*');
```

```
regexp_split_to_array
------------------------
{the, quick, brown, fox, jumps, over, the, lazy, dog}
```

(1 row)

```
SELECT foo FROM regexp_split_to_table('the quick brown fox', '\s*') AS foo;
```

<table>
<thead>
<tr>
<th>foo</th>
</tr>
</thead>
<tbody>
<tr>
<td>t</td>
</tr>
<tr>
<td>h</td>
</tr>
<tr>
<td>e</td>
</tr>
<tr>
<td>q</td>
</tr>
<tr>
<td>u</td>
</tr>
<tr>
<td>i</td>
</tr>
<tr>
<td>c</td>
</tr>
<tr>
<td>k</td>
</tr>
<tr>
<td>b</td>
</tr>
<tr>
<td>r</td>
</tr>
<tr>
<td>o</td>
</tr>
<tr>
<td>w</td>
</tr>
<tr>
<td>n</td>
</tr>
<tr>
<td>f</td>
</tr>
<tr>
<td>o</td>
</tr>
<tr>
<td>x</td>
</tr>
</tbody>
</table>

(16 rows)

As the last example demonstrates, the *regexp split* functions ignore zero-length matches that occur at the start or end of the string or immediately after a previous match. This is contrary to the strict
definition of regexp matching that is implemented by `regexp_matches`, but is usually the most convenient behavior in practice. Other software systems such as Perl use similar definitions.

### 9.7.3.1. Regular Expression Details

PostgreSQL’s regular expressions are implemented using a software package written by Henry Spencer. Much of the description of regular expressions below is copied verbatim from his manual.

Regular expressions (REs), as defined in POSIX 1003.2, come in two forms: *extended* REs or EREs (roughly those of `egrep`), and *basic* REs or BREs (roughly those of `ed`). PostgreSQL supports both forms, and also implements some extensions that are not in the POSIX standard, but have become widely used due to their availability in programming languages such as Perl and Tcl. REs using these non-POSIX extensions are called *advanced* REs or AREs in this documentation. AREs are almost an exact superset of EREs, but BREs have several notational incompatibilities (as well as being much more limited). We first describe the ARE and ERE forms, noting features that apply only to AREs, and then describe how BREs differ.

**Note:** PostgreSQL always initially presumes that a regular expression follows the ARE rules. However, the more limited ERE or BRE rules can be chosen by prepending an *embedded option* to the RE pattern, as described in Section 9.7.3.4. This can be useful for compatibility with applications that expect exactly the POSIX 1003.2 rules.

A regular expression is defined as one or more *branches*, separated by `|`. It matches anything that matches one of the branches.

A branch is zero or more *quantified atoms or constraints*, concatenated. It matches a match for the first, followed by a match for the second, etc; an empty branch matches the empty string.

A quantified atom is an *atom* possibly followed by a single *quantifier*. Without a quantifier, it matches a match for the atom. With a quantifier, it can match some number of matches of the atom. An atom can be any of the possibilities shown in Table 9-15. The possible quantifiers and their meanings are shown in Table 9-16.

A *constraint* matches an empty string, but matches only when specific conditions are met. A constraint can be used where an atom could be used, except it cannot be followed by a quantifier. The simple constraints are shown in Table 9-17; some more constraints are described later.

#### Table 9-15. Regular Expression Atoms

<table>
<thead>
<tr>
<th>Atom</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>(re)</code></td>
<td>(where <code>re</code> is any regular expression) matches a match for <code>re</code>, with the match noted for possible reporting</td>
</tr>
<tr>
<td><code>(?:re)</code></td>
<td>as above, but the match is not noted for reporting (a “non-capturing” set of parentheses) (AREs only)</td>
</tr>
<tr>
<td>.</td>
<td>matches any single character</td>
</tr>
<tr>
<td><code>[chars]</code></td>
<td>a <em>bracket expression</em>, matching any one of the <code>chars</code> (see Section 9.7.3.2 for more detail)</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Atom</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\k</td>
<td>(where (k) is a non-alphanumeric character) matches that character taken as an ordinary character, e.g., \ \ matches a backslash character</td>
</tr>
<tr>
<td>\c</td>
<td>where (c) is alphanumeric (possibly followed by other characters) is an escape, see Section 9.7.3.3 (AREs only; in EREs and BREs, this matches (c))</td>
</tr>
<tr>
<td>{</td>
<td>when followed by a character other than a digit, matches the left-brace character {; when followed by a digit, it is the beginning of a bound (see below)</td>
</tr>
<tr>
<td>x</td>
<td>where (x) is a single character with no other significance, matches that character</td>
</tr>
</tbody>
</table>

An RE cannot end with a backslash (\).  

**Note:** If you have standard_conforming_strings turned off, any backslashes you write in literal string constants will need to be doubled. See Section 4.1.2.1 for more information.

### Table 9-16. Regular Expression Quantifiers

<table>
<thead>
<tr>
<th>Quantifier</th>
<th>Matches</th>
</tr>
</thead>
<tbody>
<tr>
<td>*</td>
<td>a sequence of 0 or more matches of the atom</td>
</tr>
<tr>
<td>+</td>
<td>a sequence of 1 or more matches of the atom</td>
</tr>
<tr>
<td>?</td>
<td>a sequence of 0 or 1 matches of the atom</td>
</tr>
<tr>
<td>{m}</td>
<td>a sequence of exactly (m) matches of the atom</td>
</tr>
<tr>
<td>{m,}</td>
<td>a sequence of (m) or more matches of the atom</td>
</tr>
<tr>
<td>{m,n}</td>
<td>a sequence of (m) through (n) (inclusive) matches of the atom; (m) cannot exceed (n)</td>
</tr>
<tr>
<td>*?</td>
<td>non-greedy version of *</td>
</tr>
<tr>
<td>+?</td>
<td>non-greedy version of +</td>
</tr>
<tr>
<td>??</td>
<td>non-greedy version of ?</td>
</tr>
<tr>
<td>{m}?</td>
<td>non-greedy version of {m}</td>
</tr>
<tr>
<td>{m,}?</td>
<td>non-greedy version of {m,}</td>
</tr>
<tr>
<td>{m,n}?</td>
<td>non-greedy version of {m,n}</td>
</tr>
</tbody>
</table>

The forms using \{ . . . \} are known as **bounds**. The numbers \(m\) and \(n\) within a bound are unsigned decimal integers with permissible values from 0 to 255 inclusive.

**Non-greedy** quantifiers (available in AREs only) match the same possibilities as their corresponding normal (greedy) counterparts, but prefer the smallest number rather than the largest number of matches. See Section 9.7.3.5 for more detail.

**Note:** A quantifier cannot immediately follow another quantifier, e.g., *+* is invalid. A quantifier cannot begin an expression or subexpression or follow ^ or |.
Table 9-17. Regular Expression Constraints

<table>
<thead>
<tr>
<th>Constraint</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>^</td>
<td>matches at the beginning of the string</td>
</tr>
<tr>
<td>$</td>
<td>matches at the end of the string</td>
</tr>
<tr>
<td>(?=re)</td>
<td>positive lookahead matches at any point where a substring matching (re) begins (AREs only)</td>
</tr>
<tr>
<td>(?!re)</td>
<td>negative lookahead matches at any point where no substring matching (re) begins (AREs only)</td>
</tr>
<tr>
<td>(?&lt;re)</td>
<td>positive lookbehind matches at any point where a substring matching (re) ends (AREs only)</td>
</tr>
<tr>
<td>(?!re)</td>
<td>negative lookbehind matches at any point where no substring matching (re) ends (AREs only)</td>
</tr>
</tbody>
</table>

Lookahead and lookbehind constraints cannot contain back references (see Section 9.7.3.3), and all parentheses within them are considered non-capturing.

9.7.3.2. Bracket Expressions

A bracket expression is a list of characters enclosed in [ ]. It normally matches any single character from the list (but see below). If the list begins with ^, it matches any single character not from the rest of the list. If two characters in the list are separated by -, this is shorthand for the full range of characters between those two (inclusive) in the collating sequence, e.g., [0-9] in ASCII matches any decimal digit. It is illegal for two ranges to share an endpoint, e.g., a-c-e. Ranges are very collating-sequence-dependent, so portable programs should avoid relying on them.

To include a literal ] in the list, make it the first character (after ^, if that is used). To include a literal -, make it the first or last character, or the second endpoint of a range. To use a literal - as the first endpoint of a range, enclose it in [ . and . ] to make it a collating element (see below). With the exception of these characters, some combinations using [ (see next paragraphs), and escapes (AREs only), all other special characters lose their special significance within a bracket expression. In particular, \ is not special when following ERE or BRE rules, though it is special (as introducing an escape) in AREs.

Within a bracket expression, a collating element (a character, a multiple-character sequence that collates as if it were a single character, or a collating-sequence name for either) enclosed in [ . and . ] stands for the sequence of characters of that collating element. The sequence is treated as a single element of the bracket expression’s list. This allows a bracket expression containing a multiple-character collating element to match more than one character, e.g., if the collating sequence includes a ch collating element, then the RE [[ .ch . ]] *c matches the first five characters of chchcc.

Note: PostgreSQL currently does not support multi-character collating elements. This information describes possible future behavior.

Within a bracket expression, a collating element enclosed in [ and ] is an equivalence class, standing for the sequences of characters of all collating elements equivalent to that one, including itself. (If there are no other equivalent collating elements, the treatment is as if the enclosing delimiters were [ .

Page 220
and .) For example, if o and ^ are the members of an equivalence class, then \([=-o^{-}]\), \([=-^{-}]\), and \([o^]\) are all synonymous. An equivalence class cannot be an endpoint of a range.

Within a bracket expression, the name of a character class enclosed in \([: \text{ and } :)\) stands for the list of all characters belonging to that class. Standard character class names are: \texttt{alnum}, \texttt{alpha}, \texttt{blank}, \texttt{cntrl}, \texttt{digit}, \texttt{graph}, \texttt{lower}, \texttt{print}, \texttt{punct}, \texttt{space}, \texttt{upper}, \texttt{xdigit}. These stand for the character classes defined in \texttt{ctype}. A locale can provide others. A character class cannot be used as an endpoint of a range.

There are two special cases of bracket expressions: the bracket expressions \([[:<:]\) and \([[:>:]\) are constraints, matching empty strings at the beginning and end of a word respectively. A word is defined as a sequence of word characters that is neither preceded nor followed by word characters. A word character is an \texttt{alnum} character (as defined by \texttt{ctype}) or an underscore. This is an extension, compatible with but not specified by POSIX 1003.2, and should be used with caution in software intended to be portable to other systems. The constraint escapes described below are usually preferable; they are no more standard, but are easier to type.

### 9.7.3.3. Regular Expression Escapes

Escapes are special sequences beginning with \(\backslash\) followed by an alphanumeric character. Escapes come in several varieties: character entry, class shorthands, constraint escapes, and back references. A \(\backslash\) followed by an alphanumeric character but not constituting a valid escape is illegal in AREs. In EREs, there are no escapes: outside a bracket expression, a \(\backslash\) followed by an alphanumeric character merely stands for that character as an ordinary character, and inside a bracket expression, \(\backslash\) is an ordinary character. (The latter is the one actual incompatibility between EREs and AREs.)

**Character-entry escapes** exist to make it easier to specify non-printing and other inconvenient characters in REs. They are shown in Table 9-18.

**Class-shorthand escapes** provide shorthands for certain commonly-used character classes. They are shown in Table 9-19.

A **constraint escape** is a constraint, matching the empty string if specific conditions are met, written as an escape. They are shown in Table 9-20.

A **back reference** \((\backslash n)\) matches the same string matched by the previous parenthesized subexpression specified by the number \(n\) (see Table 9-21). For example, \(([bc])\1\) matches bb or cc but not bc or cb. The subexpression must entirely precede the back reference in the RE. Subexpressions are numbered in the order of their leading parentheses. Non-capturing parentheses do not define subexpressions.

#### Table 9-18. Regular Expression Character-entry Escapes

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\a</td>
<td>alert (bell) character, as in C</td>
</tr>
<tr>
<td>\b</td>
<td>backspace, as in C</td>
</tr>
<tr>
<td>\B</td>
<td>synonym for backslash () to help reduce the need for backslash doubling</td>
</tr>
<tr>
<td>\cX</td>
<td>(where (X) is any character) the character whose low-order 5 bits are the same as those of (X), and whose other bits are all zero</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

Escape | Description
---|---
\e | the character whose collating-sequence name is ESC, or failing that, the character with octal value 033
\f | form feed, as in C
\n | newline, as in C
\r | carriage return, as in C
\t | horizontal tab, as in C
\u{wxyz} | (where \textit{wxyz} is exactly four hexadecimal digits) the character whose hexadecimal value is \texttt{0xwxyz}
\U{stuvwxyz} | (where \textit{stuvwxyz} is exactly eight hexadecimal digits) the character whose hexadecimal value is \texttt{0xstuvwxyz}
\v | vertical tab, as in C
\x{hhh} | (where \textit{hhh} is any sequence of hexadecimal digits) the character whose hexadecimal value is \texttt{0xhhhh} (a single character no matter how many hexadecimal digits are used)
\0 | the character whose value is 0 (the null byte)
\xy | (where \textit{xy} is exactly two octal digits, and is not a \textit{back reference}) the character whose octal value is \texttt{0xy}
\xyz | (where \textit{xyz} is exactly three octal digits, and is not a \textit{back reference}) the character whose octal value is \texttt{0xyz}

Hexadecimal digits are 0-9, a-f, and A-F. Octal digits are 0-7.

Numeric character-entry escapes specifying values outside the ASCII range (0-127) have meanings dependent on the database encoding. When the encoding is UTF-8, escape values are equivalent to Unicode code points, for example \texttt{\u1234} means the character \texttt{U+1234}. For other multibyte encodings, character-entry escapes usually just specify the concatenation of the byte values for the character. If the escape value does not correspond to any legal character in the database encoding, no error will be raised, but it will never match any data.

The character-entry escapes are always taken as ordinary characters. For example, \texttt{\135} is \texttt{]} in ASCII, but \texttt{\135} does not terminate a bracket expression.

Table 9-19. Regular Expression Class-shorthand Escapes

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\d</td>
<td>[[:digit:]]</td>
</tr>
<tr>
<td>\s</td>
<td>[[:space:]]</td>
</tr>
<tr>
<td>\w</td>
<td>[[:alnum:]] (note underscore is included)</td>
</tr>
<tr>
<td>\D</td>
<td>[^[:digit:]]</td>
</tr>
<tr>
<td>\S</td>
<td>[^[:space:]]</td>
</tr>
<tr>
<td>\W</td>
<td>[^[:alnum:]] (note underscore is included)</td>
</tr>
</tbody>
</table>
Within bracket expressions, \d, \s, and \w lose their outer brackets, and \D, \S, and \W are illegal. (So, for example, \[a-c\d\] is equivalent to \[a-c[:digit:]\]. Also, \[a-c\D\], which is equivalent to \[a-c^[:digit:]\], is illegal.)

Table 9-20. Regular Expression Constraint Escapes

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\A</td>
<td>matches only at the beginning of the string (see Section 9.7.3.5 for how this differs from ^)</td>
</tr>
<tr>
<td>\m</td>
<td>matches only at the beginning of a word</td>
</tr>
<tr>
<td>\M</td>
<td>matches only at the end of a word</td>
</tr>
<tr>
<td>\y</td>
<td>matches only at the beginning or end of a word</td>
</tr>
<tr>
<td>\Y</td>
<td>matches only at a point that is not the beginning or end of a word</td>
</tr>
<tr>
<td>\Z</td>
<td>matches only at the end of the string (see Section 9.7.3.5 for how this differs from $)</td>
</tr>
</tbody>
</table>

A word is defined as in the specification of \[[:<:]\] and \[[:>:]\] above. Constraint escapes are illegal within bracket expressions.

Table 9-21. Regular Expression Back References

<table>
<thead>
<tr>
<th>Escape</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\m</td>
<td>(where \m is a nonzero digit) a back reference to the \m'th subexpression</td>
</tr>
<tr>
<td>\mnn</td>
<td>(where \m is a nonzero digit, and \nn is some more digits, and the decimal value \mnn is not greater than the number of closing capturing parentheses seen so far) a back reference to the \mnn'th subexpression</td>
</tr>
</tbody>
</table>

Note: There is an inherent ambiguity between octal character-entry escapes and back references, which is resolved by the following heuristics, as hinted at above. A leading zero always indicates an octal escape. A single non-zero digit, not followed by another digit, is always taken as a back reference. A multi-digit sequence not starting with a zero is taken as a back reference if it comes after a suitable subexpression (i.e., the number is in the legal range for a back reference), and otherwise is taken as octal.

9.7.3.4. Regular Expression Metasyntax

In addition to the main syntax described above, there are some special forms and miscellaneous syntactic facilities available.

An RE can begin with one of two special director prefixes. If an RE begins with ***:, the rest of the RE is taken as an ARE. (This normally has no effect in PostgreSQL, since REs are assumed to be AREs; but it does have an effect if ERE or BRE mode had been specified by the flags parameter to a regex function.) If an RE begins with ***=, the rest of the RE is taken to be a literal string, with all characters considered ordinary characters.
Chapter 9. Functions and Operators

An ARE can begin with embedded options: a sequence \( (?xyz) \) (where \( xyz \) is one or more alphabetic characters) specifies options affecting the rest of the RE. These options override any previously determined options — in particular, they can override the case-sensitivity behavior implied by a regex operator, or the \( flags \) parameter to a regex function. The available option letters are shown in Table 9-22. Note that these same option letters are used in the \( flags \) parameters of regex functions.

Table 9-22. ARE Embedded-option Letters

<table>
<thead>
<tr>
<th>Option</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>b</td>
<td>rest of RE is a BRE</td>
</tr>
<tr>
<td>c</td>
<td>case-sensitive matching (overrides operator type)</td>
</tr>
<tr>
<td>e</td>
<td>rest of RE is an ERE</td>
</tr>
<tr>
<td>i</td>
<td>case-insensitive matching (see Section 9.7.3.5) (overrides operator type)</td>
</tr>
<tr>
<td>m</td>
<td>historical synonym for n</td>
</tr>
<tr>
<td>n</td>
<td>newline-sensitive matching (see Section 9.7.3.5)</td>
</tr>
<tr>
<td>p</td>
<td>partial newline-sensitive matching (see Section 9.7.3.5)</td>
</tr>
<tr>
<td>q</td>
<td>rest of RE is a literal (“quoted”) string, all ordinary characters</td>
</tr>
<tr>
<td>s</td>
<td>non-newline-sensitive matching (default)</td>
</tr>
<tr>
<td>t</td>
<td>tight syntax (default; see below)</td>
</tr>
<tr>
<td>w</td>
<td>inverse partial newline-sensitive (“weird”) matching (see Section 9.7.3.5)</td>
</tr>
<tr>
<td>x</td>
<td>expanded syntax (see below)</td>
</tr>
</tbody>
</table>

Embedded options take effect at the \( ) \) terminating the sequence. They can appear only at the start of an ARE (after the **:** director if any).

In addition to the usual (tight) RE syntax, in which all characters are significant, there is an expanded syntax, available by specifying the embedded \( x \) option. In the expanded syntax, white-space characters in the RE are ignored, as are all characters between a \( # \) and the following newline (or the end of the RE). This permits paragraphing and commenting a complex RE. There are three exceptions to that basic rule:

- a white-space character or \( # \) preceded by \( \backslash \) is retained
- white space or \( # \) within a bracket expression is retained
- white space and comments cannot appear within multi-character symbols, such as \( (?:: \)

For this purpose, white-space characters are blank, tab, newline, and any character that belongs to the \( space \) character class.

Finally, in an ARE, outside bracket expressions, the sequence \( (?#ttt) \) (where \( ttt \) is any text not containing a \( ) \)) is a comment, completely ignored. Again, this is not allowed between the characters of multi-character symbols, like \( (?:: \). Such comments are more a historical artifact than a useful facility, and their use is deprecated; use the expanded syntax instead.

None of these metasyntax extensions is available if an initial **:** director has specified that the user’s input be treated as a literal string rather than as an RE.
9.7.3.5. Regular Expression Matching Rules

In the event that an RE could match more than one substring of a given string, the RE matches the one starting earliest in the string. If the RE could match more than one substring starting at that point, either the longest possible match or the shortest possible match will be taken, depending on whether the RE is **greedy** or **non-greedy**.

Whether an RE is greedy or not is determined by the following rules:

- Most atoms, and all constraints, have no greediness attribute (because they cannot match variable amounts of text anyway).
- Adding parentheses around an RE does not change its greediness.
- A quantified atom with a fixed-repetition quantifier (\( \{m\} \) or \( \{m\}? \)) has the same greediness (possibly none) as the atom itself.
- A quantified atom with other normal quantifiers (including \( \{m,n\} \) with \( m \) equal to \( n \)) is greedy (prefers longest match).
- A quantified atom with a non-greedy quantifier (including \( \{m,n\}? \) with \( m \) equal to \( n \)) is non-greedy (prefers shortest match).
- A branch — that is, an RE that has no top-level | operator — has the same greediness as the first quantified atom in it that has a greediness attribute.
- An RE consisting of two or more branches connected by the | operator is always greedy.

The above rules associate greediness attributes not only with individual quantified atoms, but with branches and entire REs that contain quantified atoms. What that means is that the matching is done in such a way that the branch, or whole RE, matches the longest or shortest possible substring as a whole. Once the length of the entire match is determined, the part of it that matches any particular subexpression is determined on the basis of the greediness attribute of that subexpression, with subexpressions starting earlier in the RE taking priority over ones starting later.

An example of what this means:

```
SELECT SUBSTRING('XY1234Z', 'Y*([0-9]{1,3})');
Result: 123
```

```
SELECT SUBSTRING('XY1234Z', 'Y*?([0-9]{1,3})');
Result: 1
```

In the first case, the RE as a whole is greedy because \( Y^* \) is greedy. It can match beginning at the \( Y \), and it matches the longest possible string starting there, i.e., \( Y123 \). The output is the parenthesized part of that, or \( 123 \). In the second case, the RE as a whole is non-greedy because \( Y^*? \) is non-greedy. It can match beginning at the \( Y \), and it matches the shortest possible string starting there, i.e., \( Y1 \). The subexpression \( [0-9]{1,3} \) is greedy but it cannot change the decision as to the overall match length; so it is forced to match just \( 1 \).

In short, when an RE contains both greedy and non-greedy subexpressions, the total match length is either as long as possible or as short as possible, according to the attribute assigned to the whole RE. The attributes assigned to the subexpressions only affect how much of that match they are allowed to “eat” relative to each other.

The quantifiers \( \{1,1\} \) and \( \{1,1\}? \) can be used to force greediness or non-greediness, respectively, on a subexpression or a whole RE. This is useful when you need the whole RE to have a greediness attribute different from what’s deduced from its elements. As an example, suppose that we are trying
to separate a string containing some digits into the digits and the parts before and after them. We might try to do that like this:

SELECT regexp_matches('abc01234xyz', '(.*)(\d+)(.*)');
Result: {abc0123,4,xyz}

That didn’t work: the first . * is greedy so it “eats” as much as it can, leaving the \d+ to match at the last possible place, the last digit. We might try to fix that by making it non-greedy:

SELECT regexp_matches('abc01234xyz', '(.*?)(\d+)(.*)');
Result: {abc,0,""}

That didn’t work either, because now the RE as a whole is non-greedy and so it ends the overall match as soon as possible. We can get what we want by forcing the RE as a whole to be greedy:

SELECT regexp_matches('abc01234xyz', '(?:(.*?)(\d+)(.*)){1,1}');
Result: {abc,01234,xyz}

Controlling the RE’s overall greediness separately from its components’ greediness allows great flexibility in handling variable-length patterns.

When deciding what is a longer or shorter match, match lengths are measured in characters, not collating elements. An empty string is considered longer than no match at all. For example: bb* matches the three middle characters of abbbc; (week|wee)(night|knighta) matches all ten characters of weeknights; when . * . * is matched against abc the parenthesized subexpression matches all three characters; and when (a*)* is matched against bc both the whole RE and the parenthesized subexpression match an empty string.

If case-independent matching is specified, the effect is much as if all case distinctions had vanished from the alphabet. When an alphabetic that exists in multiple cases appears as an ordinary character outside a bracket expression, it is effectively transformed into a bracket expression containing both cases, e.g., x becomes [xX]. When it appears inside a bracket expression, all case counterparts of it are added to the bracket expression, e.g., [x] becomes [xX] and [^x] becomes [^xX].

If newline-sensitive matching is specified, ^ will never match the newline character (so that matches will never cross newlines unless the RE explicitly arranges it) and $ will match the empty string after and before a newline respectively, in addition to matching at beginning and end of string respectively. But the ARE escapes \A and \Z continue to match beginning or end of string only.

If partial newline-sensitive matching is specified, this affects . and bracket expressions as with newline-sensitive matching, but not ^ and $.

If inverse partial newline-sensitive matching is specified, this affects ^ and $ as with newline-sensitive matching, but not . and bracket expressions. This isn’t very useful but is provided for symmetry.

9.7.3.6. Limits and Compatibility

No particular limit is imposed on the length of REs in this implementation. However, programs intended to be highly portable should not employ REs longer than 256 bytes, as a POSIX-compliant implementation can refuse to accept such REs.

The only feature of AREs that is actually incompatible with POSIX EREs is that \ does not lose its special significance inside bracket expressions. All other ARE features use syntax which is illegal or has undefined or unspecified effects in POSIX EREs; the *** syntax of directors likewise is outside the POSIX syntax for both BREs and EREs.
Many of the ARE extensions are borrowed from Perl, but some have been changed to clean them up, and a few Perl extensions are not present. Incompatibilities of note include \b, \B, the lack of special treatment for a trailing newline, the addition of complemented bracket expressions to the things affected by newline-sensitive matching, the restrictions on parentheses and back references in lookahead/lookbehind constraints, and the longest/shortest-match (rather than first-match) matching semantics.

Two significant incompatibilities exist between AREs and the ERE syntax recognized by pre-7.4 releases of PostgreSQL:

- In AREs, \ followed by an alphanumeric character is either an escape or an error, while in previous releases, it was just another way of writing the alphanumeric. This should not be much of a problem because there was no reason to write such a sequence in earlier releases.
- In AREs, \ remains a special character within [], so a literal \ within a bracket expression must be written \\.

9.7.3.7. Basic Regular Expressions

BREs differ from EREs in several respects. In BREs, |, +, and ? are ordinary characters and there is no equivalent for their functionality. The delimiters for bounds are \\ and \}, with { and } by themselves ordinary characters. The parentheses for nested subexpressions are \( and \), with ( and ) by themselves ordinary characters. ^ is an ordinary character except at the beginning of the RE or the beginning of a parenthesized subexpression, $ is an ordinary character except at the end of the RE or the end of a parenthesized subexpression, and * is an ordinary character if it appears at the beginning of the RE or the beginning of a parenthesized subexpression (after a possible leading ^). Finally, single-digit back references are available, and \< and \> are synonyms for [[<:]] and [[>:]] respectively; no other escapes are available in BREs.

9.8. Data Type Formatting Functions

The PostgreSQL formatting functions provide a powerful set of tools for converting various data types (date/time, integer, floating point, numeric) to formatted strings and for converting from formatted strings to specific data types. Table 9-23 lists them. These functions all follow a common calling convention: the first argument is the value to be formatted and the second argument is a template that defines the output or input format.

<table>
<thead>
<tr>
<th>Table 9-23. Formatting Functions</th>
</tr>
</thead>
<tbody>
<tr>
<td>Function</td>
</tr>
<tr>
<td>---------</td>
</tr>
<tr>
<td>to_char(timestamp, text)</td>
</tr>
<tr>
<td>to_char(interval, text)</td>
</tr>
<tr>
<td>to_char(int, text)</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>to_char(double precision, text)</td>
<td>text</td>
<td>convert real/double precision to string</td>
<td>to_char(125.8::real, '999D9')</td>
</tr>
<tr>
<td>to_char(numeric, text)</td>
<td>text</td>
<td>convert numeric to string</td>
<td>to_char(-125.8, '999D9S')</td>
</tr>
<tr>
<td>to_date(text, text)</td>
<td>date</td>
<td>convert string to date</td>
<td>to_date('05 Dec 2000', 'DD Mon YYYY')</td>
</tr>
<tr>
<td>to_number(text, text)</td>
<td>numeric</td>
<td>convert string to numeric</td>
<td>to_number('12,454.8-', '99G9999D9S')</td>
</tr>
<tr>
<td>to_timestamp(text, text)</td>
<td>timestamp with time zone</td>
<td>convert string to time stamp</td>
<td>to_timestamp('05 Dec 2000', 'DD Mon YYYY')</td>
</tr>
</tbody>
</table>

**Note:** There is also a single-argument `to_timestamp` function; see Table 9-30.

In a to_char output template string, there are certain patterns that are recognized and replaced with appropriately-formatted data based on the given value. Any text that is not a template pattern is simply copied verbatim. Similarly, in an input template string (for the other functions), template patterns identify the values to be supplied by the input data string.

Table 9-24 shows the template patterns available for formatting date and time values.

**Table 9-24. Template Patterns for Date/Time Formatting**

<table>
<thead>
<tr>
<th>Pattern</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>HH</td>
<td>hour of day (01-12)</td>
</tr>
<tr>
<td>HH12</td>
<td>hour of day (01-12)</td>
</tr>
<tr>
<td>HH24</td>
<td>hour of day (00-23)</td>
</tr>
<tr>
<td>MI</td>
<td>minute (00-59)</td>
</tr>
<tr>
<td>SS</td>
<td>second (00-59)</td>
</tr>
<tr>
<td>MS</td>
<td>millisecond (000-999)</td>
</tr>
<tr>
<td>US</td>
<td>microsecond (000000-999999)</td>
</tr>
<tr>
<td>SSSS</td>
<td>seconds past midnight (0-86399)</td>
</tr>
<tr>
<td>AM, am, PM or pm</td>
<td>meridiem indicator (without periods)</td>
</tr>
<tr>
<td>A.M., a.m., P.M. or p.m.</td>
<td>meridiem indicator (with periods)</td>
</tr>
<tr>
<td>Y, YYYY</td>
<td>year (4 or more digits) with comma</td>
</tr>
<tr>
<td>YYY</td>
<td>year (4 or more digits)</td>
</tr>
<tr>
<td>YY</td>
<td>last 3 digits of year</td>
</tr>
<tr>
<td>Y</td>
<td>last 2 digits of year</td>
</tr>
<tr>
<td>Y</td>
<td>last digit of year</td>
</tr>
<tr>
<td>YYYY</td>
<td>ISO 8601 week-numbering year (4 or more digits)</td>
</tr>
<tr>
<td>IYYYY</td>
<td>ISO 8601 week-numbering year</td>
</tr>
<tr>
<td>yyyy</td>
<td>last 3 digits of ISO 8601 week-numbering year</td>
</tr>
<tr>
<td>iy</td>
<td>last 2 digits of ISO 8601 week-numbering year</td>
</tr>
<tr>
<td>I</td>
<td>last digit of ISO 8601 week-numbering year</td>
</tr>
<tr>
<td>BC, bc, AD or ad</td>
<td>era indicator (without periods)</td>
</tr>
<tr>
<td>Pattern</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>B.C., b.c., A.D. or a.d.</td>
<td>era indicator (with periods)</td>
</tr>
<tr>
<td>MONTH</td>
<td>full upper case month name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>Month</td>
<td>full capitalized month name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>month</td>
<td>full lower case month name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>MON</td>
<td>abbreviated upper case month name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>Mon</td>
<td>abbreviated capitalized month name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>mon</td>
<td>abbreviated lower case month name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>MM</td>
<td>month number (01-12)</td>
</tr>
<tr>
<td>DAY</td>
<td>full upper case day name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>Day</td>
<td>full capitalized day name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>day</td>
<td>full lower case day name (blank-padded to 9 chars)</td>
</tr>
<tr>
<td>DY</td>
<td>abbreviated upper case day name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>Dy</td>
<td>abbreviated capitalized day name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>dy</td>
<td>abbreviated lower case day name (3 chars in English, localized lengths vary)</td>
</tr>
<tr>
<td>DDD</td>
<td>day of year (001-366)</td>
</tr>
<tr>
<td>IDDD</td>
<td>day of ISO 8601 week-numbering year (001-371; day 1 of the year is Monday of the first ISO week)</td>
</tr>
<tr>
<td>DD</td>
<td>day of month (01-31)</td>
</tr>
<tr>
<td>D</td>
<td>day of the week, Sunday (1) to Saturday (7)</td>
</tr>
<tr>
<td>ID</td>
<td>ISO 8601 day of the week, Monday (1) to Sunday (7)</td>
</tr>
<tr>
<td>W</td>
<td>week of month (1-5) (the first week starts on the first day of the month)</td>
</tr>
<tr>
<td>WW</td>
<td>week number of year (1-53) (the first week starts on the first day of the year)</td>
</tr>
<tr>
<td>IN</td>
<td>week number of ISO 8601 week-numbering year (01-53; the first Thursday of the year is in week 1)</td>
</tr>
<tr>
<td>CC</td>
<td>century (2 digits) (the twenty-first century starts on 2001-01-01)</td>
</tr>
<tr>
<td>J</td>
<td>Julian Day (integer days since November 24, 4714 BC at midnight UTC)</td>
</tr>
</tbody>
</table>
Modifiers can be applied to any template pattern to alter its behavior. For example, `FMMonth` is the `Month` pattern with the `FM` modifier. Table 9-25 shows the modifier patterns for date/time formatting.

Table 9-25. Template Pattern Modifiers for Date/Time Formatting

<table>
<thead>
<tr>
<th>Modifier</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>FM</strong> prefix</td>
<td>fill mode (suppress leading zeroes and padding blanks)</td>
<td><code>FMMonth</code></td>
</tr>
<tr>
<td><strong>TH</strong> suffix</td>
<td>upper case ordinal number suffix</td>
<td><code>DDTH</code>, e.g., <code>12TH</code></td>
</tr>
<tr>
<td><strong>th</strong> suffix</td>
<td>lower case ordinal number suffix</td>
<td><code>DDth</code>, e.g., <code>12th</code></td>
</tr>
<tr>
<td><strong>FX</strong> prefix</td>
<td>fixed format global option (see usage notes)</td>
<td><code>FX Month DD Day</code></td>
</tr>
<tr>
<td><strong>TM</strong> prefix</td>
<td>translation mode (print localized day and month names based on <code>lc_time</code>)</td>
<td><code>TMMonth</code></td>
</tr>
<tr>
<td><strong>SP</strong> suffix</td>
<td>spell mode (not implemented)</td>
<td><code>DDSP</code></td>
</tr>
</tbody>
</table>

Usage notes for date/time formatting:

- **FM** suppresses leading zeroes and trailing blanks that would otherwise be added to make the output of a pattern be fixed-width. In PostgreSQL, `FM` modifies only the next specification, while in Oracle `FM` affects all subsequent specifications, and repeated `FM` modifiers toggle fill mode on and off.
- **TM** does not include trailing blanks. `to_timestamp` and `to_date` ignore the `TM` modifier.
- `to_timestamp` and `to_date` skip multiple blank spaces in the input string unless the `FX` option is used. For example, `to_timestamp('2000 JUN', 'YYYY MM')` works, but `to_timestamp('2000 JUN', 'FXYYYY MON')` returns an error because `to_timestamp` expects one space only. `FX` must be specified as the first item in the template.
- `to_timestamp` and `to_date` exist to handle input formats that cannot be converted by simple casting. These functions interpret input liberally, with minimal error checking. While they produce valid output, the conversion can yield unexpected results. For example, input to these functions is not restricted by normal ranges, thus `to_date('20096040','YYYYMMDD')` returns `2014-01-17`
rather than causing an error. Casting does not have this behavior.

- Ordinary text is allowed in `to_char` templates and will be output literally. You can put a substring in double quotes to force it to be interpreted as literal text even if it contains pattern key words. For example, in `"Hello Year "YYYY"`, the `YYYY` will be replaced by the year data, but the single `Y` in `Year` will not be. In `to_date`, `to_number`, and `to_timestamp`, double-quoted strings skip the number of input characters contained in the string, e.g. "xx" skips two input characters.

- If you want to have a double quote in the output you must precede it with a backslash, for example `'"YYYY Month"'`.

- If the year format specification is less than four digits, e.g. `YYY`, and the supplied year is less than four digits, the year will be adjusted to be nearest to the year 2020, e.g. `95` becomes `1995`.

- The `YYYY` conversion from string to `timestamp` or `date` has a restriction when processing years with more than 4 digits. You must use some non-digit character or template after `YYYY`, otherwise the year is always interpreted as 4 digits. For example (with the year `20000`): `to_date('200001131', 'YYYYMMD')` will be interpreted as a 4-digit year; instead use a non-digit separator after the year, like `to_date('20000-1131', 'YYYY-MMDD')` or `to_date('20000Nov31', 'YYYYMonDD')`.

- In conversions from string to `timestamp` or `date`, the `CC` (century) field is ignored if there is a `YYY`, `YYYY` or `Y,YYY` field. If `CC` is used with `YY` or `Y` then the year is computed as the year in the specified century. If the century is specified but the year is not, the first year of the century is assumed.

- An ISO 8601 week-numbering date (as distinct from a Gregorian date) can be specified to `to_timestamp` and `to_date` in one of two ways:
  - Year, week number, and weekday: for example `to_date('2006-42-4', 'IYYY-IW-ID')` returns the date `2006-10-19`. If you omit the weekday it is assumed to be 1 (Monday).
  - Year and day of year: for example `to_date('2006-291', 'IYYYY-IDDD')` also returns `2006-10-19`.

Attempting to enter a date using a mixture of ISO 8601 week-numbering fields and Gregorian date fields is nonsensical, and will cause an error. In the context of an ISO 8601 week-numbering year, the concept of a “month” or “day of month” has no meaning. In the context of a Gregorian year, the ISO week has no meaning.

```
Caution
While `to_date` will reject a mixture of Gregorian and ISO week-numbering date fields, `to_char` will not, since output format specifications like `YYYY-MM-DD` (IYYYY-IDDD) can be useful. But avoid writing something like `IYYY-MM-DD`; that would yield surprising results near the start of the year. (See Section 9.9.1 for more information.)
```

- In a conversion from string to `timestamp`, millisecond (MS) or microsecond (US) values are used as the seconds digits after the decimal point. For example `to_timestamp('12:3', 'SS:MS')` is not 3 milliseconds, but 300, because the conversion counts it as `12 + 0.3` seconds. This means for the format `SS:MS`, the input values `12:3`, `12:30`, and `12:300` specify the same number of milliseconds. To get three milliseconds, one must use `12:003`, which the conversion counts as `12 + 0.003 = 12.003` seconds.
Chapter 9. Functions and Operators

Here is a more complex example: `to_timestamp('15:12:02.020.001230', 'HH24:MI:SS:MS.US')` is 15 hours, 12 minutes, and 2 seconds + 20 milliseconds + 1230 microseconds = 2.021230 seconds.

- `to_char(..., 'ID')`’s day of the week numbering matches the `extract(isodow from ...)` function, but `to_char(..., 'D')`’s does not match `extract(dow from ...)`’s day numbering.

- `to_char(interval)` formats `HH` and `HH12` as shown on a 12-hour clock, i.e. zero hours and 36 hours output as 12, while `HH24` outputs the full hour value, which can exceed 23 for intervals.

Table 9-26 shows the template patterns available for formatting numeric values.

<table>
<thead>
<tr>
<th>Pattern</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>9</td>
<td>digit position (can be dropped if insignificant)</td>
</tr>
<tr>
<td>0</td>
<td>digit position (will not be dropped, even if insignificant)</td>
</tr>
<tr>
<td>, (period)</td>
<td>decimal point</td>
</tr>
<tr>
<td>, (comma)</td>
<td>group (thousands) separator</td>
</tr>
<tr>
<td>PR</td>
<td>negative value in angle brackets</td>
</tr>
<tr>
<td>$</td>
<td>sign anchored to number (uses locale)</td>
</tr>
<tr>
<td>L</td>
<td>currency symbol (uses locale)</td>
</tr>
<tr>
<td>D</td>
<td>decimal point (uses locale)</td>
</tr>
<tr>
<td>G</td>
<td>group separator (uses locale)</td>
</tr>
<tr>
<td>MI</td>
<td>minus sign in specified position (if number &lt; 0)</td>
</tr>
<tr>
<td>PL</td>
<td>plus sign in specified position (if number &gt; 0)</td>
</tr>
<tr>
<td>SG</td>
<td>plus/minus sign in specified position</td>
</tr>
<tr>
<td>RN</td>
<td>Roman numeral (input between 1 and 3999)</td>
</tr>
<tr>
<td>TH or th</td>
<td>ordinal number suffix</td>
</tr>
<tr>
<td>V</td>
<td>shift specified number of digits (see notes)</td>
</tr>
<tr>
<td>EEEE</td>
<td>exponent for scientific notation</td>
</tr>
</tbody>
</table>

Usage notes for numeric formatting:

- 0 specifies a digit position that will always be printed, even if it contains a leading/trailing zero. 9 also specifies a digit position, but if it is a leading zero then it will be replaced by a space, while if it is a trailing zero and fill mode is specified then it will be deleted. (For `to_number()`, these two pattern characters are equivalent.)

- The pattern characters $, L, D,$ and $G$ represent the sign, currency symbol, decimal point, and thousands separator characters defined by the current locale (see `lc_monetary` and `lc_numeric`). The pattern characters period and comma represent those exact characters, with the meanings of decimal point and thousands separator, regardless of locale.

- If no explicit provision is made for a sign in `to_char()`’s pattern, one column will be reserved for the sign, and it will be anchored to (appear just left of) the number. If $S$ appears just left of some 9’s, it will likewise be anchored to the number.
A sign formatted using `SG`, `PL`, or `MI` is not anchored to the number; for example, `to_char(-12, 'MI9999')` produces `' 12'` but `to_char(-12, 'S9999')` produces `' -12'`. (The Oracle implementation does not allow the use of `MI` before 9, but rather requires that 9 precede `MI`.)

- `TH` does not convert values less than zero and does not convert fractional numbers.
- `PL`, `SG`, and `TH` are PostgreSQL extensions.
- `V` with `to_char` multiplies the input values by $10^n$, where $n$ is the number of digits following `V`. `V` with `to_number` divides in a similar manner. `to_char` and `to_number` do not support the use of `V` combined with a decimal point (e.g., `.9999` is not allowed).
- `EEEE` (scientific notation) cannot be used in combination with any of the other formatting patterns or modifiers other than digit and decimal point patterns, and must be at the end of the format string (e.g., `9.999999` is a valid pattern).

Certain modifiers can be applied to any template pattern to alter its behavior. For example, `FM99.99` is the `99.99` pattern with the `FM` modifier. Table 9-27 shows the modifier patterns for numeric formatting.

### Table 9-27. Template Pattern Modifiers for Numeric Formatting

<table>
<thead>
<tr>
<th>Modifier</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>FM</td>
<td>prefix</td>
<td><code>FM99.99</code></td>
</tr>
<tr>
<td>TH</td>
<td>suffix</td>
<td><code>999TH</code></td>
</tr>
<tr>
<td>th</td>
<td>suffix</td>
<td><code>999th</code></td>
</tr>
</tbody>
</table>

Table 9-28 shows some examples of the use of the `to_char` function.

### Table 9-28. `to_char` Examples

<table>
<thead>
<tr>
<th>Expression</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>to_char(current_timestamp, 'Day, DD HH12:MI:SS')</code></td>
<td>'Tuesday, 06 05:39:18'</td>
</tr>
<tr>
<td><code>to_char(current_timestamp, 'FMDay, FMDD HH12:MI:SS')</code></td>
<td>'Tuesday, 6 05:39:18'</td>
</tr>
<tr>
<td><code>to_char(-0.1, '99.99')</code></td>
<td>'-.10'</td>
</tr>
<tr>
<td><code>to_char(-0.1, 'FM9.99')</code></td>
<td>'-.1'</td>
</tr>
<tr>
<td><code>to_char(-0.1, 'FM90.99')</code></td>
<td>'-0.1'</td>
</tr>
<tr>
<td><code>to_char(0.1, '0.9')</code></td>
<td>'0.1'</td>
</tr>
<tr>
<td><code>to_char(12, '9990999.9')</code></td>
<td>'0012.0'</td>
</tr>
<tr>
<td><code>to_char(12, 'FM9990999.9')</code></td>
<td>'0012.'</td>
</tr>
<tr>
<td><code>to_char(485, '999')</code></td>
<td>'485'</td>
</tr>
<tr>
<td><code>to_char(-485, '999')</code></td>
<td>'-485'</td>
</tr>
<tr>
<td><code>to_char(485, '9 9 9')</code></td>
<td>'4 8 5'</td>
</tr>
<tr>
<td><code>to_char(1485, '9,999')</code></td>
<td>'1,485'</td>
</tr>
<tr>
<td><code>to_char(1485, '9G999')</code></td>
<td>'1 485'</td>
</tr>
</tbody>
</table>
9.9. Date/Time Functions and Operators

Table 9-30 shows the available functions for date/time value processing, with details appearing in the following subsections. Table 9-29 illustrates the behaviors of the basic arithmetic operators (+, *, etc.). For formatting functions, refer to Section 9.8. You should be familiar with the background information on date/time data types from Section 8.5.

All the functions and operators described below that take time or timestamp inputs actually come in two variants: one that takes time with time zone or timestamp with time zone, and one that takes time without time zone or timestamp without time zone. For brevity, these variants are not shown separately. Also, the + and * operators come in commutative pairs (for example both date + integer and integer + date); we show only one of each such pair.

Table 9-29. Date/Time Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>to_char(148.5, '999.999')</td>
<td>'148.500'</td>
<td></td>
</tr>
<tr>
<td>to_char(148.5, 'FM999.999')</td>
<td>'148.5'</td>
<td></td>
</tr>
<tr>
<td>to_char(148.5, 'FM999.990')</td>
<td>'148.500'</td>
<td></td>
</tr>
<tr>
<td>to_char(148.5, '999D999')</td>
<td>'148,500'</td>
<td></td>
</tr>
<tr>
<td>to_char(3148.5, '9G999D999')</td>
<td>'3 148,500'</td>
<td></td>
</tr>
<tr>
<td>to_char(-485, '999S')</td>
<td>'-485-'</td>
<td></td>
</tr>
<tr>
<td>to_char(-485, '999MI')</td>
<td>'-485-'</td>
<td></td>
</tr>
<tr>
<td>to_char(485, '999MI')</td>
<td>'+485'</td>
<td></td>
</tr>
<tr>
<td>to_char(485, 'PL999')</td>
<td>'+485'</td>
<td></td>
</tr>
<tr>
<td>to_char(485, 'SG999')</td>
<td>'-485'</td>
<td></td>
</tr>
<tr>
<td>to_char(-485, 'SG999')</td>
<td>'4-85'</td>
<td></td>
</tr>
<tr>
<td>to_char(-485, '9SG99')</td>
<td>'&lt;-85&gt;'</td>
<td></td>
</tr>
<tr>
<td>to_char(485, '1999')</td>
<td>'D 485'</td>
<td></td>
</tr>
<tr>
<td>to_char(485, 'RN')</td>
<td>' CDLXXXV'</td>
<td></td>
</tr>
<tr>
<td>to_char(485, 'FMRN')</td>
<td>'CDLXXXV'</td>
<td></td>
</tr>
<tr>
<td>to_char(5.2, 'FMRN')</td>
<td>'V'</td>
<td></td>
</tr>
<tr>
<td>to_char(482, '999th')</td>
<td>'482nd'</td>
<td></td>
</tr>
<tr>
<td>to_char(485, '&quot;Good number:&quot;999')</td>
<td>'Good number: 485'</td>
<td></td>
</tr>
<tr>
<td>to_char(485.8, '&quot;Pre:&quot;999&quot; Post:&quot; .999')</td>
<td>'Pre: 485 Post: .800'</td>
<td></td>
</tr>
<tr>
<td>to_char(12, '&quot;99V999')</td>
<td>'12000'</td>
<td></td>
</tr>
<tr>
<td>to_char(12.4, '&quot;99V999')</td>
<td>'12400'</td>
<td></td>
</tr>
<tr>
<td>to_char(12.45, '&quot;99V9')</td>
<td>'125'</td>
<td></td>
</tr>
<tr>
<td>to_char(0.0004859, '&quot;9.99EEEE')</td>
<td>'4.86e-04'</td>
<td></td>
</tr>
<tr>
<td>Operator</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>---------</td>
<td>---------</td>
<td>-----------------</td>
</tr>
<tr>
<td>+</td>
<td>date '2001-09-28' + integer '7'</td>
<td>date '2001-10-05'</td>
</tr>
<tr>
<td>+</td>
<td>date '2001-09-28' + interval '1 hour'</td>
<td>timestamp '2001-09-28 01:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>date '2001-09-28' + time '03:00'</td>
<td>timestamp '2001-09-28 03:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>interval '1 day' + interval '1 hour'</td>
<td>interval '1 day 01:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>timestamp '2001-09-28 01:00' + interval '23 hours'</td>
<td>timestamp '2001-09-29 00:00:00'</td>
</tr>
<tr>
<td>+</td>
<td>time '01:00' + interval '3 hours'</td>
<td>time '04:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>- interval '23 hours'</td>
<td>interval '-23:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>date '2001-10-01' - date '2001-09-28'</td>
<td>integer '3' (days)</td>
</tr>
<tr>
<td>-</td>
<td>date '2001-10-01' - integer '7'</td>
<td>date '2001-09-24'</td>
</tr>
<tr>
<td>-</td>
<td>date '2001-09-28' - interval '1 hour'</td>
<td>timestamp '2001-09-27 23:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>time '05:00' - time '03:00'</td>
<td>interval '02:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>time '05:00' - interval '2 hours'</td>
<td>time '03:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>timestamp '2001-09-28 23:00' - interval '23 hours'</td>
<td>timestamp '2001-09-28 00:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>interval '1 day' - interval '1 hour'</td>
<td>interval '1 day -01:00:00'</td>
</tr>
<tr>
<td>-</td>
<td>timestamp '2001-09-29 03:00' - timestamp '2001-09-27 12:00'</td>
<td>interval '1 day 15:00:00'</td>
</tr>
<tr>
<td>*</td>
<td>900 * interval '1 second'</td>
<td>interval '00:15:00'</td>
</tr>
<tr>
<td>*</td>
<td>21 * interval '1 day'</td>
<td>interval '21 days'</td>
</tr>
<tr>
<td>*</td>
<td>double precision '3.5' * interval '1 hour'</td>
<td>interval '03:30:00'</td>
</tr>
<tr>
<td>/</td>
<td>interval '1 hour' / double precision '1.5'</td>
<td>interval '00:40:00'</td>
</tr>
</tbody>
</table>

Table 9-30. Date/Time Functions
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>age(timestamp, timestamp)</td>
<td>interval</td>
<td>Subtract arguments, producing a &quot;symbolic&quot; result that uses years and months, rather than just days</td>
<td>age(timestamp '2001-04-10',</td>
<td>43 years 9 mons 27 days</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>timestamp '1957-06-13')</td>
<td></td>
</tr>
<tr>
<td>age(timestamp)</td>
<td>interval</td>
<td>Subtract from current_date (at midnight)</td>
<td>age(timestamp '1957-06-13')</td>
<td>43 years 8 mons 3 days</td>
</tr>
<tr>
<td>clock_timestamp()</td>
<td>timestamp with</td>
<td>Current date and time (changes during statement execution); see</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>time zone</td>
<td>Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>current_date</td>
<td>date</td>
<td>Current date; see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>current_time</td>
<td>time with time</td>
<td>Current time of day; see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>zone</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>current_timestamp</td>
<td>timestamp with</td>
<td>Current date and time (start of current transaction); see</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td>time zone</td>
<td>Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>date_part(text, timestamp)</td>
<td>double precision</td>
<td>Get subfield (equivalent to extract); see Section 9.9.1</td>
<td>date_part('hour', '2001-02-16</td>
<td>20:38:40')</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>20:38:40')</td>
<td></td>
</tr>
<tr>
<td>date_part(text, interval)</td>
<td>double precision</td>
<td>Get subfield (equivalent to extract); see Section 9.9.1</td>
<td>date_part('month', '2001-02-16</td>
<td>20:38:40')</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>20:38:40')</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>year '2' months)</td>
<td></td>
</tr>
<tr>
<td>date_trunc(text, timestamp)</td>
<td>timestamp</td>
<td>Truncate to specified precision; see also Section 9.9.2</td>
<td>date_trunc('hour', '2001-02-16</td>
<td>20:00:00</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>20:38:40')</td>
<td></td>
</tr>
<tr>
<td>date_trunc(text, interval)</td>
<td>interval</td>
<td>Truncate to specified precision; see also Section 9.9.2</td>
<td>date_trunc('hour', '2 days 3</td>
<td>03:00:00</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>hours 40 minutes')</td>
<td></td>
</tr>
<tr>
<td>extract(field from timestamp)</td>
<td>double precision</td>
<td>Get subfield; see Section 9.9.1</td>
<td>extract('hour from timestamp</td>
<td>20</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>'2001-02-16 20:38:40')</td>
<td></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>-------------</td>
<td>---------</td>
<td>--------</td>
</tr>
<tr>
<td><code>extract(field from interval)</code></td>
<td><code>double precision</code></td>
<td>Get subfield; see Section 9.9.1</td>
<td><code>extract(month from interval '2 years 3 months')</code></td>
<td>3</td>
</tr>
<tr>
<td><code>isfinite(date)</code></td>
<td><code>boolean</code></td>
<td>Test for finite date (not +/-infinity)</td>
<td><code>isfinite(date '2001-02-16')</code></td>
<td>true</td>
</tr>
<tr>
<td><code>isfinite(timestamp)</code></td>
<td><code>boolean</code></td>
<td>Test for finite timestamp (not +/-infinity)</td>
<td><code>isfinite(timestamp '2001-02-16 21:28:30')</code></td>
<td></td>
</tr>
<tr>
<td><code>isfinite(interval)</code></td>
<td><code>boolean</code></td>
<td>Test for finite interval</td>
<td><code>isfinite(interval '4 hours')</code></td>
<td>true</td>
</tr>
<tr>
<td><code>justify_days(interval)</code></td>
<td><code>interval</code></td>
<td>Adjust interval so 30-day time periods are represented as months</td>
<td><code>justify_days(interval '35 days')</code></td>
<td></td>
</tr>
<tr>
<td><code>justify_hours(interval)</code></td>
<td><code>interval</code></td>
<td>Adjust interval so 24-hour time periods are represented as days</td>
<td><code>justify_hours(interval '27 hours')</code></td>
<td>03:00:00</td>
</tr>
<tr>
<td><code>justify_interval(interval)</code></td>
<td><code>interval</code></td>
<td>Adjust interval using <code>justify_days</code> and <code>justify_hours</code>, with additional sign adjustments</td>
<td><code>justify_interval(interval '1 mon -1 hour')</code></td>
<td>23:00:00</td>
</tr>
<tr>
<td><code>localtime</code></td>
<td><code>time</code></td>
<td>Current time of day; see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>localtimestamp</code></td>
<td><code>timestamp</code></td>
<td>Current date and time (start of current transaction); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>make_date(year int, month int, day int)</code></td>
<td><code>date</code></td>
<td>Create date from year, month and day fields</td>
<td><code>make_date(2013, 7, 15)</code></td>
<td>2013-07-15</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>------------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
<td>-----------------------</td>
<td>--------------</td>
</tr>
<tr>
<td>make_interval</td>
<td>interval</td>
<td>Create interval from years, months, weeks, days, hours, minutes and seconds</td>
<td>make_interval(days ( \geq 10 ))</td>
<td>10 days</td>
</tr>
<tr>
<td>make_time</td>
<td>time</td>
<td>Create time from hour, minute and seconds fields</td>
<td>make_time(8, 15, 23.5)</td>
<td>08:15:23.5</td>
</tr>
<tr>
<td>make_timestamp</td>
<td>timestamp</td>
<td>Create timestamp from year, month, day, hour, minute and seconds fields</td>
<td>make_timestamp(2013, 7, 15, 8, 15, 23.5)</td>
<td>2013-07-15 08:15:23.5</td>
</tr>
<tr>
<td>make_timestampz</td>
<td>timestamp</td>
<td>Create timestamp with time zone from year, month, day, hour, minute and</td>
<td>make_timestampz(2013, 7, 15, 8, 15, 23.5)</td>
<td>2013-07-15 08:15:23.5+01</td>
</tr>
<tr>
<td></td>
<td></td>
<td>seconds fields; if timezone is not specified, the current time zone is used</td>
<td></td>
<td></td>
</tr>
<tr>
<td>now()</td>
<td>timestamp</td>
<td>Current date and time (start of current transaction); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td>statement_timestamp</td>
<td>timestamp</td>
<td>Current date and time (start of current statement); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>timeofday()</code></td>
<td>text</td>
<td>Current date and time (like <code>clock_timestamp</code>, but as a text string); see 47</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>transaction_timestamp()</code></td>
<td>timestamp</td>
<td>Current date and time (start of current transaction); see Section 9.9.4</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>to_timestamp(double precision)</code></td>
<td>timestamp</td>
<td>Convert Unix epoch (seconds since 1970-01-01 00:00:00+00) to timestamp</td>
<td><code>to_timestamp(1284352323)</code></td>
<td><code>2010-09-13 04:32:03+00</code></td>
</tr>
</tbody>
</table>

In addition to these functions, the SQL `OVERLAPS` operator is supported:

```
(start1, end1) OVERLAPS (start2, end2)
(start1, length1) OVERLAPS (start2, length2)
```

This expression yields true when two time periods (defined by their endpoints) overlap, false when they do not overlap. The endpoints can be specified as pairs of dates, times, or time stamps; or as a date, time, or time stamp followed by an interval. When a pair of values is provided, either the start or the end can be written first; `OVERLAPS` automatically takes the earlier value of the pair as the start. Each time period is considered to represent the half-open interval `start <= time < end`, unless `start` and `end` are equal in which case it represents that single time instant. This means for instance that two time periods with only an endpoint in common do not overlap.

```
SELECT (DATE '2001-02-16', DATE '2001-12-21') OVERLAPS
     (DATE '2001-10-30', DATE '2002-10-30');
Result: true
```

```
SELECT (DATE '2001-02-16', INTERVAL '100 days') OVERLAPS
     (DATE '2001-10-30', DATE '2002-10-30');
Result: false
```

```
SELECT (DATE '2001-10-29', DATE '2001-10-30') OVERLAPS
     (DATE '2001-10-30', DATE '2001-10-31');
Result: false
```

```
SELECT (DATE '2001-10-30', DATE '2001-10-30') OVERLAPS
     (DATE '2001-10-30', DATE '2001-10-31');
Result: true
```

When adding an `interval` value to (or subtracting an `interval` value from) a `timestamp` with `time zone` value, the days component advances or decrements the date of the `timestamp` with `time zone` by the indicated number of days. Across daylight saving time changes (when the session time zone is set to a time zone that recognizes DST), this means `interval` ‘1 day’ does not necessarily equal interval ‘24 hours’. For example, with the session time zone set to CST7CDT, `timestamp` with time zone ‘2005-04-02 12:00-07’ + interval ‘1 day’ will produce `timestamp` with time zone ‘2005-04-03 12:00-06’, while adding `interval` ‘24 hours’ to the same initial `timestamp` with time zone produces `timestamp` with time zone ‘2005-04-03 12:00-07’.
zone ‘2005-04-03 13:00:06’, as there is a change in daylight saving time at 2005-04-03 02:00 in time zone CST7CDT.

Note there can be ambiguity in the months field returned by age because different months have different numbers of days. PostgreSQL’s approach uses the month from the earlier of the two dates when calculating partial months. For example, age(‘2004-06-01’, ’2004-04-30’) uses April to yield 1 mon 1 day, while using May would yield 1 mon 2 days because May has 31 days, while April has only 30.

Subtraction of dates and timestamps can also be complex. One conceptually simple way to perform subtraction is to convert each value to a number of seconds using \texttt{EXTRACT(EPOCH FROM ...)}, then subtract the results; this produces the number of \textit{seconds} between the two values. This will adjust for the number of days in each month, timezone changes, and daylight saving time adjustments. Subtraction of date or timestamp values with the “-” operator returns the number of days (24-hours) and hours/minutes/seconds between the values, making the same adjustments. The \texttt{age} function returns years, months, days, and hours/minutes/seconds, performing field-by-field subtraction and then adjusting for negative field values. The following queries illustrate the differences in these approaches. The sample results were produced with \texttt{timezone = ‘US/Eastern’}; there is a daylight saving time change between the two dates used:

\begin{verbatim}
SELECT EXTRACT(EPOCH FROM timestamptz '2013-07-01 12:00:00') -
  EXTRACT(EPOCH FROM timestamptz '2013-03-01 12:00:00');
Result: 10537200
SELECT (EXTRACT(EPOCH FROM timestamptz '2013-07-01 12:00:00') -
  EXTRACT(EPOCH FROM timestamptz '2013-03-01 12:00:00'))
 / 60 / 60 / 24;
Result: 121.958333333333
SELECT timestamptz '2013-07-01 12:00:00' - timestamptz '2013-03-01 12:00:00';
Result: 121 days 23:00:00
SELECT age(timestamptz '2013-07-01 12:00:00', timestamptz '2013-03-01 12:00:00');
Result: 4 mons
\end{verbatim}

9.9.1. \texttt{EXTRACT}, \texttt{date_part}

\texttt{EXTRACT(field FROM source)}

The \texttt{extract} function retrieves subfields such as year or hour from date/time values. \texttt{source} must be a value expression of type \texttt{timestamp}, \texttt{time}, or \texttt{interval}. (Expressions of type \texttt{date} are cast to \texttt{timestamp} and can therefore be used as well.) \texttt{field} is an identifier or string that selects what field to extract from the source value. The \texttt{extract} function returns values of type \texttt{double precision}. The following are valid field names:

- \texttt{century}
  
The century
  
  \begin{verbatim}
  SELECT EXTRACT(CENTURY FROM TIMESTAMP '2000-12-16 12:21:13');
  Result: 20
  SELECT EXTRACT(CENTURY FROM TIMESTAMP '2001-02-16 20:38:40');
  Result: 21
  \end{verbatim}

  The first century starts at 0001-01-01 00:00:00 AD, although they did not know it at the time. This definition applies to all Gregorian calendar countries. There is no century number 0, you go from -1 century to 1 century. If you disagree with this, please write your complaint to: Pope, Cathedral Saint-Peter of Roma, Vatican.

240
Chapter 9. Functions and Operators

day
For timestamp values, the day (of the month) field (1 - 31); for interval values, the number of days

SELECT EXTRACT(DAY FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 16

SELECT EXTRACT(DAY FROM INTERVAL '40 days 1 minute');
Result: 40
decade
The year field divided by 10

SELECT EXTRACT(DECADE FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 200
dow
The day of the week as Sunday (0) to Saturday (6)

SELECT EXTRACT(DOW FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 5

Note that extract’s day of the week numbering differs from that of the to_char(..., ‘D’) function.
doy
The day of the year (1 - 365/366)

SELECT EXTRACT(DOY FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 47
epoch
For timestamp with time zone values, the number of seconds since 1970-01-01 00:00:00 UTC (can be negative); for date and timestamp values, the number of seconds since 1970-01-01 00:00:00 local time; for interval values, the total number of seconds in the interval

SELECT EXTRACT(EPOCH FROM TIMESTAMP WITH TIME ZONE '2001-02-16 20:38:40.12-08');
Result: 982384720.12

SELECT EXTRACT(EPOCH FROM INTERVAL '5 days 3 hours');
Result: 442800

You can convert an epoch value back to a time stamp with to_timestamp:

SELECT to_timestamp(982384720.12);
Result: 2001-02-17 04:38:40.12+00

hour
The hour field (0 - 23)

SELECT EXTRACT(HOUR FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 20

isodow
The day of the week as Monday (1) to Sunday (7)

SELECT EXTRACT(ISODOW FROM TIMESTAMP '2001-02-18 20:38:40');
Result: 7

This is identical to dow except for Sunday. This matches the ISO 8601 day of the week numbering.
isoyear

The ISO 8601 week-numbering year that the date falls in (not applicable to intervals)

```sql
SELECT EXTRACT(ISOYEAR FROM DATE '2006-01-01');
Result: 2005

SELECT EXTRACT(ISOYEAR FROM DATE '2006-01-02');
Result: 2006
```

Each ISO 8601 week-numbering year begins with the Monday of the week containing the 4th of
January, so in early January or late December the ISO year may be different from the Gregorian
year. See the `week` field for more information.

This field is not available in PostgreSQL releases prior to 8.3.

microseconds

The seconds field, including fractional parts, multiplied by 1 000 000; note that this includes full
seconds

```sql
SELECT EXTRACT(MICROSECONDS FROM TIME '17:12:28.5');
Result: 28500000
```

millennium

The millennium

```sql
SELECT EXTRACT(MILLENNIUM FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 3
```

Years in the 1900s are in the second millennium. The third millennium started January 1, 2001.

milliseconds

The seconds field, including fractional parts, multiplied by 1000. Note that this includes full
seconds.

```sql
SELECT EXTRACT(MILLISECONDS FROM TIME '17:12:28.5');
Result: 28500
```

minute

The minutes field (0 - 59)

```sql
SELECT EXTRACT(MINUTE FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 38
```

month

For `timestamp` values, the number of the month within the year (1 - 12); for `interval` values,
the number of months, modulo 12 (0 - 11)

```sql
SELECT EXTRACT(MONTH FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 2
```

```sql
SELECT EXTRACT(MONTH FROM INTERVAL '2 years 3 months');
Result: 3
```

```sql
SELECT EXTRACT(MONTH FROM INTERVAL '2 years 13 months');
Result: 1
```

quarter

The quarter of the year (1 - 4) that the date is in

```sql
SELECT EXTRACT(QUARTER FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 1
```
Chapter 9. Functions and Operators

second

The seconds field, including fractional parts (0 - 59½)

```
SELECT EXTRACT(SECOND FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 40
```

```
SELECT EXTRACT(SECOND FROM TIME '17:12:28.5');
Result: 28.5
```

timezone

The time zone offset from UTC, measured in seconds. Positive values correspond to time zones east of UTC, negative values to zones west of UTC. (Technically, PostgreSQL does not use UTC because leap seconds are not handled.)

timezone_hour

The hour component of the time zone offset

timezone_minute

The minute component of the time zone offset

week

The number of the ISO 8601 week-numbering week of the year. By definition, ISO weeks start on Mondays and the first week of a year contains January 4 of that year. In other words, the first Thursday of a year is in week 1 of that year.

In the ISO week-numbering system, it is possible for early-January dates to be part of the 52nd or 53rd week of the previous year, and for late-December dates to be part of the first week of the next year. For example, 2005-01-01 is part of the 53rd week of year 2004, and 2006-01-01 is part of the 52nd week of year 2005, while 2012-12-31 is part of the first week of 2013. It’s recommended to use the isoyear field together with week to get consistent results.

```
SELECT EXTRACT(WEEK FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 7
```

year

The year field. Keep in mind there is no 0 AD, so subtracting BC years from AD years should be done with care.

```
SELECT EXTRACT(YEAR FROM TIMESTAMP '2001-02-16 20:38:40');
Result: 2001
```

Note: When the input value is +/-Infinity, extract returns +/-Infinity for monotonically-increasing fields (epoch, julian, year, isoyear, decade, century, and millennium). For other fields, NULL is returned. PostgreSQL versions before 9.6 returned zero for all cases of infinite input.

The `extract` function is primarily intended for computational processing. For formatting date/time values for display, see Section 9.8.

The `date_part` function is modeled on the traditional Ingres equivalent to the SQL-standard function `extract`:

```
date_part('field', source)
```

60 if leap seconds are implemented by the operating system
Note that here the field parameter needs to be a string value, not a name. The valid field names for date_part are the same as for extract.

SELECT date_part('day', TIMESTAMP '2001-02-16 20:38:40');
Result: 16

SELECT date_part('hour', INTERVAL '4 hours 3 minutes');
Result: 4

### 9.9.2. date_trunc

The function date_trunc is conceptually similar to the trunc function for numbers.

```
date_trunc('field', source)
```

source is a value expression of type timestamp or interval. (Values of type date and time are cast automatically to timestamp or interval, respectively.) field selects to which precision to truncate the input value. The return value is of type timestamp or interval with all fields that are less significant than the selected one set to zero (or one, for day and month).

Valid values for field are:

- microseconds
- milliseconds
- second
- minute
- hour
- day
- week
- month
- quarter
- year
- decade
- century
- millennium

Examples:

```
SELECT date_trunc('hour', TIMESTAMP '2001-02-16 20:38:40');
Result: 2001-02-16 20:00:00

SELECT date_trunc('year', TIMESTAMP '2001-02-16 20:38:40');
Result: 2001-01-01 00:00:00
```

### 9.9.3. AT TIME ZONE

The AT TIME ZONE converts time stamp without time zone to/from time stamp with time zone, and time values to different time zones. Table 9-31 shows its variants.
Table 9-31. AT TIME ZONE Variants

<table>
<thead>
<tr>
<th>Expression</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>timestamp without time zone AT TIME ZONE zone</td>
<td>timestamp with time zone</td>
<td>Treat given time stamp without time zone as located in the specified time zone</td>
</tr>
<tr>
<td>timestamp with time zone AT TIME ZONE zone</td>
<td>timestamp without time zone</td>
<td>Convert given time stamp with time zone to the new time zone, with no time zone designation</td>
</tr>
<tr>
<td>time with time zone AT TIME ZONE zone</td>
<td>time with time zone</td>
<td>Convert given time with time zone to the new time zone</td>
</tr>
</tbody>
</table>

In these expressions, the desired time zone \textit{zone} can be specified either as a text string (e.g., 'America/Los_Angeles') or as an interval (e.g., \texttt{INTERVAL '-08:00')}). In the text case, a time zone name can be specified in any of the ways described in Section 8.5.3.

Examples (assuming the local time zone is America/Los_Angeles):

\begin{verbatim}
SELECT TIMESTAMP '2001-02-16 20:38:40' AT TIME ZONE 'America/Denver';
Result: 2001-02-16 19:38:40-08

SELECT TIMESTAMP WITH TIME ZONE '2001-02-16 20:38:40-05' AT TIME ZONE 'America/Denver';
Result: 2001-02-16 18:38:40

SELECT TIMESTAMP '2001-02-16 20:38:40-05' AT TIME ZONE 'Asia/Tokyo' AT TIME ZONE 'America/Chicago';
Result: 2001-02-16 05:38:40
\end{verbatim}

The first example adds a time zone to a value that lacks it, and displays the value using the current TimeZone setting. The second example shifts the time stamp with time zone value to the specified time zone, and returns the value without a time zone. This allows storage and display of values different from the current TimeZone setting. The third example converts Tokyo time to Chicago time.

Converting time values to other time zones uses the currently active time zone rules since no date is supplied.

The function \texttt{timezone(zone, timestamp)} is equivalent to the SQL-conforming construct \texttt{timestamp AT TIME ZONE zone}.

9.9.4. Current Date/Time

PostgreSQL provides a number of functions that return values related to the current date and time. These SQL-standard functions all return values based on the start time of the current transaction:

\begin{verbatim}
CURRENT_DATE
CURRENT_TIME
CURRENT_TIMESTAMP
CURRENT_TIME(precision)
CURRENT_TIMESTAMP(precision)
LOCALTIME
LOCALTIMESTAMP
LOCALTIME(precision)
LOCALTIMESTAMP(precision)
\end{verbatim}
CURRENT_TIME and CURRENT_TIMESTAMP deliver values with time zone; LOCALTIME and LOCALTIMESTAMP deliver values without time zone.

CURRENT_TIME, CURRENT_TIMESTAMP, LOCALTIME, and LOCALTIMESTAMP can optionally take a precision parameter, which causes the result to be rounded to that many fractional digits in the seconds field. Without a precision parameter, the result is given to the full available precision.

Some examples:

```
SELECT CURRENT_TIME;
Result: 14:39:53.662522-05
```

```
SELECT CURRENT_DATE;
Result: 2001-12-23
```

```
SELECT CURRENT_TIMESTAMP;
Result: 2001-12-23 14:39:53.662522-05
```

```
SELECT CURRENT_TIMESTAMP(2);
Result: 2001-12-23 14:39:53.66-05
```

```
SELECT LOCALTIMESTAMP;
Result: 2001-12-23 14:39:53.662522
```

Since these functions return the start time of the current transaction, their values do not change during the transaction. This is considered a feature: the intent is to allow a single transaction to have a consistent notion of the “current” time, so that multiple modifications within the same transaction bear the same timestamp.

**Note:** Other database systems might advance these values more frequently.

PostgreSQL also provides functions that return the start time of the current statement, as well as the actual current time at the instant the function is called. The complete list of non-SQL-standard time functions is:

- `transaction_timestamp()`
- `statement_timestamp()`
- `clock_timestamp()`
- `timeofday()`
- `now()`

`transaction_timestamp()` is equivalent to CURRENT_TIMESTAMP, but is named to clearly reflect what it returns. `statement_timestamp()` returns the start time of the current statement (more specifically, the time of receipt of the latest command message from the client). `statement_timestamp()` and `transaction_timestamp()` return the same value during the first command of a transaction, but might differ during subsequent commands. `clock_timestamp()` returns the actual current time, and therefore its value changes even within a single SQL command. `timeofday()` is a historical PostgreSQL function. Like `clock_timestamp()`, it returns the actual current time, but as a formatted text string rather than a timestamp with time zone value. `now()` is a traditional PostgreSQL equivalent to `transaction_timestamp()`.
All the date/time data types also accept the special literal value `now` to specify the current date and time (again, interpreted as the transaction start time). Thus, the following three all return the same result:

```
SELECT CURRENT_TIMESTAMP;
SELECT now();
SELECT TIMESTAMP 'now'; -- incorrect for use with DEFAULT
```

Tip: You do not want to use the third form when specifying a `DEFAULT` clause while creating a table. The system will convert `now` to a `timestamp` as soon as the constant is parsed, so that when the default value is needed, the time of the table creation would be used! The first two forms will not be evaluated until the default value is used, because they are function calls. Thus they will give the desired behavior of defaulting to the time of row insertion.

### 9.9.5. Delaying Execution

The following functions are available to delay execution of the server process:

- `pg_sleep(seconds)`
- `pg_sleep_for(interval)`
- `pg_sleep_until(timestamp with time zone)`

`pg_sleep` makes the current session’s process sleep until `seconds` seconds have elapsed. `seconds` is a value of type `double precision`, so fractional-second delays can be specified. `pg_sleep_for` is a convenience function for larger sleep times specified as an `interval`. `pg_sleep_until` is a convenience function for when a specific wake-up time is desired. For example:

```
SELECT pg_sleep(1.5);
SELECT pg_sleep_for('5 minutes');
SELECT pg_sleep_until('tomorrow 03:00');
```

Note: The effective resolution of the sleep interval is platform-specific; 0.01 seconds is a common value. The sleep delay will be at least as long as specified. It might be longer depending on factors such as server load. In particular, `pg_sleep_until` is not guaranteed to wake up exactly at the specified time, but it will not wake up any earlier.

Warning

Make sure that your session does not hold more locks than necessary when calling `pg_sleep` or its variants. Otherwise other sessions might have to wait for your sleeping process, slowing down the entire system.
9.10. Enum Support Functions

For enum types (described in Section 8.7), there are several functions that allow cleaner programming without hard-coding particular values of an enum type. These are listed in Table 9-32. The examples assume an enum type created as:

```
CREATE TYPE rainbow AS ENUM ('red', 'orange', 'yellow', 'green', 'blue', 'purple');
```

Table 9-32. Enum Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>enum_first(anyenum)</code></td>
<td>Returns the first value of the input enum type</td>
<td><code>enum_first(null::rainbow)</code></td>
<td></td>
</tr>
<tr>
<td><code>enum_last(anyenum)</code></td>
<td>Returns the last value of the input enum type</td>
<td><code>enum_last(null::rainbow)</code></td>
<td></td>
</tr>
<tr>
<td><code>enum_range(anyenum)</code></td>
<td>Returns all values of the input enum type in an ordered array</td>
<td><code>enum_range(null::rainbow)</code></td>
<td></td>
</tr>
<tr>
<td><code>enum_range(anyenum, anyenum)</code></td>
<td>Returns the range between the two given enum values, as an ordered array. The values must be from the same enum type. If the first parameter is null, the result will start with the first value of the enum type. If the second parameter is null, the result will end with the last value of the enum type.</td>
<td><code>enum_range('orange'::rainbow, 'green'::rainbow)</code></td>
<td><code>{orange,yellow,green}</code></td>
</tr>
</tbody>
</table>

Notice that except for the two-argument form of `enum_range`, these functions disregard the specific value passed to them; they care only about its declared data type. Either null or a specific value of the type can be passed, with the same result. It is more common to apply these functions to a table column or function argument than to a hardwired type name as suggested by the examples.

9.11. Geometric Functions and Operators

The geometric types `point`, `box`, `lseg`, `line`, `path`, `polygon`, and `circle` have a large set of native support functions and operators, shown in Table 9-33, Table 9-34, and Table 9-35.
Chapter 9. Functions and Operators

Caution

Note that the “same as” operator, `~=` represents the usual notion of equality for the point, box, polygon, and circle types. Some of these types also have an `=` operator, but `=` compares for equal areas only. The other scalar comparison operators (`<` and so on) likewise compare areas for these types.

Table 9-33. Geometric Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>+</td>
<td>Translation</td>
<td>box '((0,0),(1,1))' + point '(2.0,0)'</td>
</tr>
<tr>
<td>-</td>
<td>Translation</td>
<td>box '((0,0),(1,1))' - point '(2.0,0)'</td>
</tr>
<tr>
<td>*</td>
<td>Scaling/rotation</td>
<td>box '((0,0),(1,1))' * point '(2.0,0)'</td>
</tr>
<tr>
<td>/</td>
<td>Scaling/rotation</td>
<td>box '((0,0),(2,2))' / point '(2.0,0)'</td>
</tr>
<tr>
<td>#</td>
<td>Point or box of intersection</td>
<td>box '((1,-1),(-1,1))' # box '((1,1),(-2,-2))'</td>
</tr>
<tr>
<td>#</td>
<td>Number of points in path or polygon</td>
<td># path '((1,0),(0,1),(-1,0))'</td>
</tr>
<tr>
<td>@-@</td>
<td>Length or circumference</td>
<td>@-@ path '((0,0),(1,0))'</td>
</tr>
<tr>
<td>@@</td>
<td>Center</td>
<td>@@ circle '((0,0),10)'</td>
</tr>
<tr>
<td>##</td>
<td>Closest point to first operand on second operand</td>
<td>point '(0,0)' ## lseg '((2,0),(0,2))'</td>
</tr>
<tr>
<td>&lt;--&gt;</td>
<td>Distance between</td>
<td>circle '((0,0),1)' &lt;--&gt; circle '((5,0),1)'</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>Overlaps? (One point in common makes this true.)</td>
<td>box '((0,0),(1,1))' &amp;&amp; box '((0,0),(2,2))'</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>Is strictly left of?</td>
<td>circle '((0,0),1)' &lt;&lt; circle '((5,0),1)'</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>Is strictly right of?</td>
<td>circle '((5,0),1)' &gt;&gt; circle '((0,0),1)'</td>
</tr>
<tr>
<td>&amp;&lt;</td>
<td>Does not extend to the right of?</td>
<td>box '((0,0),(1,1))' &amp;&lt; box '((0,0),(2,2))'</td>
</tr>
<tr>
<td>&amp;&gt;</td>
<td>Does not extend to the left of?</td>
<td>box '((0,0),(3,3))' &amp;&gt; box '((0,0),(2,2))'</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td></td>
<td>Is strictly below?</td>
</tr>
<tr>
<td></td>
<td>&gt;&gt;</td>
<td>Is strictly above?</td>
</tr>
<tr>
<td>&amp;&lt;</td>
<td></td>
<td>Does not extend above?</td>
</tr>
<tr>
<td></td>
<td>&amp;&gt;</td>
<td>Does not extend below?</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

#### Operator

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>&lt;^</code></td>
<td>Is below (allows touching)?</td>
<td>circle <code>((0,0),1)</code> <code>&lt;^</code> circle <code>((0,5),1)</code></td>
</tr>
<tr>
<td><code>&gt;^</code></td>
<td>Is above (allows touching)?</td>
<td>circle <code>((0,5),1)</code> <code>&gt;^</code> circle <code>((0,0),1)</code></td>
</tr>
<tr>
<td><code>?#</code></td>
<td>Intersects?</td>
<td>lseg <code>((-1,0),(1,0))</code> <code>?#</code> box <code>((-2,-2),(2,2))</code></td>
</tr>
<tr>
<td><code>?-</code></td>
<td>Is horizontal?</td>
<td><code>- lseg </code>((-1,0),(1,0))`</td>
</tr>
<tr>
<td><code>?-</code></td>
<td>Are horizontally aligned?</td>
<td>point <code>'(1,0)'</code> <code>?-</code> point <code>'(0,0)'</code></td>
</tr>
<tr>
<td>`?</td>
<td>`</td>
<td>Is vertical?</td>
</tr>
<tr>
<td>`?</td>
<td>`</td>
<td>Are vertically aligned?</td>
</tr>
<tr>
<td>`?-</td>
<td>`</td>
<td>Is perpendicular?</td>
</tr>
<tr>
<td>`?</td>
<td></td>
<td>`</td>
</tr>
<tr>
<td><code>@&gt;</code></td>
<td>Contains?</td>
<td>circle <code> '((0,0),2)'</code> <code>@&gt;</code> point <code>'(1,1)'</code></td>
</tr>
<tr>
<td><code>&lt;@</code></td>
<td>Contained in or on?</td>
<td>point <code>'(1,1)'</code> <code>&lt;@</code> circle <code> '((0,0),2)'</code></td>
</tr>
<tr>
<td><code>~=</code></td>
<td>Same as?</td>
<td>polygon <code>'((0,0),(1,1))'</code>~=<code>polygon</code> '((1,1),(0,0))'`</td>
</tr>
</tbody>
</table>

**Note:** Before PostgreSQL 8.2, the containment operators `@>` and `<@` were respectively called `~` and `@`. These names are still available, but are deprecated and will eventually be removed.

#### Table 9-34. Geometric Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>area(object)</td>
<td>double precision</td>
<td>area</td>
<td>area(box <code> '((0,0),(1,1))'</code> )</td>
</tr>
<tr>
<td>center(object)</td>
<td>point</td>
<td>center</td>
<td>center(box <code> '((0,0),(1,2))'</code> )</td>
</tr>
<tr>
<td>diameter(circle)</td>
<td>double precision</td>
<td>diameter of circle</td>
<td>diameter(circle <code> '((0,0),2.0)'</code> )</td>
</tr>
<tr>
<td>height(box)</td>
<td>double precision</td>
<td>vertical size of box</td>
<td>height(box <code> '((0,0),(1,1))'</code> )</td>
</tr>
</tbody>
</table>
### Table 9-35. Geometric Type Conversion Functions

<table>
<thead>
<tr>
<th><strong>Function</strong></th>
<th><strong>Return Type</strong></th>
<th><strong>Description</strong></th>
<th><strong>Example</strong></th>
</tr>
</thead>
<tbody>
<tr>
<td><code>box(circle)</code></td>
<td>box</td>
<td>circle to box</td>
<td><code>box(circle '((0,0),2.0)')</code></td>
</tr>
<tr>
<td><code>box(point)</code></td>
<td>box</td>
<td>point to empty box</td>
<td><code>box(point '(0,0)')</code></td>
</tr>
<tr>
<td><code>box(point, point)</code></td>
<td>box</td>
<td>points to box</td>
<td><code>box(point '(0,0)', point '(1,1)')</code></td>
</tr>
<tr>
<td><code>box(polygon)</code></td>
<td>box</td>
<td>polygon to box</td>
<td><code>box(polygon '((0,0),(1,1),(2,0))')</code></td>
</tr>
<tr>
<td><code>bound_box(box, box)</code></td>
<td>box</td>
<td>boxes to bounding box</td>
<td><code>bound_box(box '((0,0),(1,1))', box '((3,3),(4,4))')</code></td>
</tr>
<tr>
<td><code>circle(box)</code></td>
<td>circle</td>
<td>box to circle</td>
<td><code>circle(box '((0,0),(1,1))')</code></td>
</tr>
<tr>
<td><code>circle(point, double precision)</code></td>
<td>circle</td>
<td>center and radius to circle</td>
<td><code>circle(point '(0,0)', 2.0)</code>)</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
</tr>
<tr>
<td>-------------------</td>
<td>-------------</td>
<td>------------------------------</td>
<td>----------------------------------------------</td>
</tr>
<tr>
<td>circle(polygon)</td>
<td>circle</td>
<td>polygon to circle</td>
<td><code>circle(polygon '((0,0),(1,1),(2,0))')')</code></td>
</tr>
<tr>
<td>line(point, point)</td>
<td>line</td>
<td>points to line</td>
<td><code>line(point '(-1,0)', point '(1,0)')')</code></td>
</tr>
<tr>
<td>lseg(box)</td>
<td>lseg</td>
<td>box diagonal to line segment</td>
<td><code>lseg(box '((-1,0),(1,0))')')</code></td>
</tr>
<tr>
<td>lseg(point, point)</td>
<td>lseg</td>
<td>points to line segment</td>
<td><code>lseg(point '(-1,0)', point '(1,0)')')</code></td>
</tr>
<tr>
<td>path(polygon)</td>
<td>path</td>
<td>polygon to path</td>
<td><code>path(polygon '((0,0),(1,1),(2,0))')')</code></td>
</tr>
<tr>
<td>point(double precision, double precision)</td>
<td>point</td>
<td>construct point</td>
<td><code>point(23.4, -44.5)</code></td>
</tr>
<tr>
<td>point(box)</td>
<td>point</td>
<td>center of box</td>
<td><code>point(box '((-1,0),(1,0))')')</code></td>
</tr>
<tr>
<td>point(circle)</td>
<td>point</td>
<td>center of circle</td>
<td><code>point(circle '((0,0),2.0)')')</code></td>
</tr>
<tr>
<td>point(lseg)</td>
<td>point</td>
<td>center of line segment</td>
<td><code>point(lseg '((-1,0),(1,0))')')</code></td>
</tr>
<tr>
<td>point(polygon)</td>
<td>point</td>
<td>center of polygon</td>
<td><code>point(polygon '((0,0),(1,1),(2,0))')')</code></td>
</tr>
<tr>
<td>polygon(box)</td>
<td>polygon</td>
<td>box to 4-point polygon</td>
<td><code>polygon(box '((-1,0),(1,1))')')</code></td>
</tr>
<tr>
<td>polygon(circle)</td>
<td>polygon</td>
<td>circle to 12-point polygon</td>
<td><code>polygon(circle '((0,0),(1,1))')')</code></td>
</tr>
<tr>
<td>polygon(npts, circle)</td>
<td>polygon</td>
<td>circle to npts-point polygon</td>
<td><code>polygon(12, circle '((0,0),(1,1))')')</code></td>
</tr>
<tr>
<td>polygon(path)</td>
<td>polygon</td>
<td>path to polygon</td>
<td><code>polygon(path '((-1,0),(1,1))')')</code></td>
</tr>
</tbody>
</table>

It is possible to access the two component numbers of a point as though the point were an array with indexes 0 and 1. For example, if `t.p` is a point column then `SELECT p[0] FROM t` retrieves the X coordinate and `UPDATE t SET p[1] = ...` changes the Y coordinate. In the same way, a value of type box or lseg can be treated as an array of two point values.

The area function works for the types box, circle, and path. The area function only works on the path data type if the points in the path are non-intersecting. For example, the path `''((0,0),(0,1),(2,1),(2,2),(1,2),(1,0),(0,0))''::PATH` will not work; however, the following visually identical path `''((0,0),(0,1),(1,1),(1,2),(2,2),(2,1),(1,1),(1,0),(0,0))''::PATH` will work. If the concept of an intersecting versus non-intersecting path is confusing, draw both of the above
paths side by side on a piece of graph paper.

9.12. Network Address Functions and Operators

Table 9-36 shows the operators available for the \texttt{cidr} and \texttt{inet} types. The operators \texttt{<<}, \texttt{<<=}, \texttt{>>}, \texttt{>>=}, and \texttt{&&} test for subnet inclusion. They consider only the network parts of the two addresses (ignoring any host part) and determine whether one network is identical to or a subnet of the other.

Table 9-36. \texttt{cidr} and \texttt{inet} Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;</td>
<td>is less than</td>
<td>\texttt{inet ’192.168.1.5’ &lt; inet ’192.168.1.6’}</td>
</tr>
<tr>
<td>&lt;=</td>
<td>is less than or equal</td>
<td>\texttt{inet ’192.168.1.5’ &lt;= inet ’192.168.1.5’}</td>
</tr>
<tr>
<td>=</td>
<td>equals</td>
<td>\texttt{inet ’192.168.1.5’ = inet ’192.168.1.5’}</td>
</tr>
<tr>
<td>&gt;=</td>
<td>is greater or equal</td>
<td>\texttt{inet ’192.168.1.5’ &gt;= inet ’192.168.1.5’}</td>
</tr>
<tr>
<td>&gt;</td>
<td>is greater than</td>
<td>\texttt{inet ’192.168.1.5’ &gt; inet ’192.168.1.4’}</td>
</tr>
<tr>
<td>&lt;&gt;</td>
<td>is not equal</td>
<td>\texttt{inet ’192.168.1.5’ &lt;&gt; inet ’192.168.1.6’}</td>
</tr>
<tr>
<td>&lt;&lt;=</td>
<td>is contained by</td>
<td>\texttt{inet ’192.168.1.5’ &lt;&lt;= inet ’192.168.1/24’}</td>
</tr>
<tr>
<td>&lt;&lt;=</td>
<td>is contained by or equals</td>
<td>\texttt{inet ’192.168.1/24’ &lt;&lt;= inet ’192.168.1/24’}</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>contains</td>
<td>\texttt{inet ’192.168.1/24’ &gt;&gt; inet ’192.168.1/24’}</td>
</tr>
<tr>
<td>&gt;&gt;=</td>
<td>contains or equals</td>
<td>\texttt{inet ’192.168.1/24’ &gt;&gt;= inet ’192.168.1/24’}</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>contains or is contained by</td>
<td>\texttt{inet ’192.168.1/24’ &amp;&amp; inet ’192.168.1/24’}</td>
</tr>
<tr>
<td>~</td>
<td>bitwise NOT</td>
<td>\texttt{~ inet ’192.168.1.6’}</td>
</tr>
<tr>
<td>&amp;</td>
<td>bitwise AND</td>
<td>\texttt{inet ’192.168.1.6’ &amp; inet ’0.0.0.255’}</td>
</tr>
<tr>
<td></td>
<td></td>
<td>bitwise OR</td>
</tr>
<tr>
<td>+</td>
<td>addition</td>
<td>\texttt{inet ’192.168.1.6’ + 25}</td>
</tr>
<tr>
<td>-</td>
<td>subtraction</td>
<td>\texttt{inet ’192.168.1.43’ - 36}</td>
</tr>
<tr>
<td>-</td>
<td>subtraction</td>
<td>\texttt{inet ’192.168.1.19’}</td>
</tr>
</tbody>
</table>

Table 9-37 shows the functions available for use with the \texttt{cidr} and \texttt{inet} types. The \texttt{abbrev}, \texttt{host}, and \texttt{text} functions are primarily intended to offer alternative display formats.
Table 9-37. cidr and inet Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>abbrev(inet)</td>
<td>text</td>
<td>abbreviated display format as text</td>
<td>abbrev(inet '10.1.0.0/16')</td>
<td>10.1.0.0/16</td>
</tr>
<tr>
<td>abbrev(cidr)</td>
<td>text</td>
<td>abbreviated display format as text</td>
<td>abbrev(cidr '10.1.0.0/16')</td>
<td>10.1/16</td>
</tr>
<tr>
<td>broadcast(inet)</td>
<td>inet</td>
<td>broadcast address for network</td>
<td>broadcast('192.168.168.2455/24')</td>
<td></td>
</tr>
<tr>
<td>family(inet)</td>
<td>int</td>
<td>extract family of address; 4 for IPv4, 6 for IPv6</td>
<td>family('::1')</td>
<td>6</td>
</tr>
<tr>
<td>host(inet)</td>
<td>text</td>
<td>extract IP address as text</td>
<td>host('192.168.192.168.1.5')</td>
<td></td>
</tr>
<tr>
<td>hostmask(inet)</td>
<td>inet</td>
<td>construct host mask for network</td>
<td>hostmask('192.168.023.20/30')</td>
<td></td>
</tr>
<tr>
<td>masklen(inet)</td>
<td>int</td>
<td>extract netmask length</td>
<td>masklen('192.168.41.5/24')</td>
<td></td>
</tr>
<tr>
<td>netmask(inet)</td>
<td>inet</td>
<td>construct netmask for network</td>
<td>netmask('192.168.255.255/255.255/255.0')</td>
<td></td>
</tr>
<tr>
<td>network(inet)</td>
<td>cidr</td>
<td>extract network part of address</td>
<td>network('192.168.49.5/24.10/24')</td>
<td></td>
</tr>
<tr>
<td>set_masklen(inet, int)</td>
<td>inet</td>
<td>set netmask length for inet value</td>
<td>set_masklen('192.168.292.6868.5/5/16')</td>
<td>192.168.292.6868.5/5/16</td>
</tr>
<tr>
<td>set_masklen(cidr, int)</td>
<td>cidr</td>
<td>set netmask length for cidr value</td>
<td>set_masklen('192.168.292.6868.0/5/16 cidr, 16)</td>
<td>192.168.292.6868.0/5/16 cidr, 16</td>
</tr>
<tr>
<td>text(inet)</td>
<td>text</td>
<td>extract IP address and netmask length as text</td>
<td>text(inet '192.168.1.5')</td>
<td>192.168.1.5/32</td>
</tr>
<tr>
<td>inet_same_family(inet, inet)</td>
<td>boolean</td>
<td>are the addresses from the same family?</td>
<td>inet_same_family('::1'0/24', '::1')</td>
<td></td>
</tr>
<tr>
<td>inet_merge(inet, inet)</td>
<td>cidr</td>
<td>the smallest network which includes both of the given networks</td>
<td>inet_merge('192.168.168.0/24', '192.168.2/5/24')</td>
<td></td>
</tr>
</tbody>
</table>

Any cidr value can be cast to inet implicitly or explicitly; therefore, the functions shown above as operating on inet also work on cidr values. (Where there are separate functions for inet and cidr, it is because the behavior should be different for the two cases.) Also, it is permitted to cast an inet value to cidr. When this is done, any bits to the right of the netmask are silently zeroed to create a valid cidr value. In addition, you can cast a text value to inet or cidr using normal casting syntax: for example, inet(expression) or colname::cidr.

Table 9-38 shows the functions available for use with the macaddr type. The function trunc(macaddr)
returns a MAC address with the last 3 bytes set to zero. This can be used to associate the remaining prefix with a manufacturer.

**Table 9-38. macaddr Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>trunc(macaddr)</td>
<td>macaddr</td>
<td>set last 3 bytes to zero</td>
<td>trunc(macaddr '12:34:56:78:90:ab')</td>
<td>12:34:56:00:00:00</td>
</tr>
</tbody>
</table>

The macaddr type also supports the standard relational operators (>, <, =, etc.) for lexicographical ordering, and the bitwise arithmetic operators (~, & and |) for NOT, AND and OR.

### 9.13. Text Search Functions and Operators

Table 9-39, Table 9-40 and Table 9-41 summarize the functions and operators that are provided for full text searching. See Chapter 12 for a detailed explanation of PostgreSQL’s text search facility.

**Table 9-39. Text Search Operators**

<table>
<thead>
<tr>
<th>Operator</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>@@</td>
<td>boolean</td>
<td>tsvector matches tsquery ?</td>
<td>to_tsvector('fat cats ate rats') @@ to_tsquery('cat &amp; rat')</td>
<td>t</td>
</tr>
<tr>
<td>@@@</td>
<td>boolean</td>
<td>deprecated synonym for @@</td>
<td>to_tsvector('fat cats ate rats') @@ to_tsquery('cat &amp; rat')</td>
<td>t</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>tsvector</td>
<td>concatenate tsvectors</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>tsquery</td>
<td>AND tsqueries together</td>
<td>'fat</td>
<td>rat':tsquery &amp;&amp; 'cat':tsquery</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>tsquery</td>
<td>OR tsqueries together</td>
</tr>
<tr>
<td>!!</td>
<td>tsquery</td>
<td>negate a tsquery</td>
<td>!! 'cat':tsquery</td>
<td>'!cat'</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;~&gt;</td>
<td>tsquery</td>
<td>tsquery</td>
<td>to_tsquery('fat' &lt;~&gt; 'rat')</td>
<td>&lt;<del>&gt; 'fat' &lt;</del>&gt; 'rat'</td>
</tr>
<tr>
<td>@&gt;</td>
<td>boolean</td>
<td>tsquery</td>
<td>'cat': to_tsquery('fat') @&gt; 'cat' &amp; 'rat': to_tsquery('rat')</td>
<td></td>
</tr>
<tr>
<td>&lt;@</td>
<td>boolean</td>
<td>tsquery</td>
<td>'cat': to_tsquery('fat') &lt;@ 'cat' &amp; 'rat': to_tsquery('rat')</td>
<td></td>
</tr>
</tbody>
</table>

#### Note:
The `tsquery` containment operators consider only the lexemes listed in the two queries, ignoring the combining operators.

In addition to the operators shown in the table, the ordinary B-tree comparison operators (=, <, etc) are defined for types `tsvector` and `tsquery`. These are not very useful for text searching but allow, for example, unique indexes to be built on columns of these types.

### Table 9-40. Text Search Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>array_to_tsvector</td>
<td>tsvector</td>
<td>convert array of lexemes to tsvector</td>
<td>array_to_tsvector('{fat,cat,rat}'::text[])</td>
<td>'fat' 'cat' 'rat'</td>
</tr>
<tr>
<td>get_current_ts_config()</td>
<td>regconfig</td>
<td>get default text search configuration</td>
<td>get_current_ts_config()</td>
<td>english</td>
</tr>
<tr>
<td>length(tsvector)</td>
<td>integer</td>
<td>number of lexemes in tsvector</td>
<td>length('fat:2,4 cat:3 rat:5A':tsvector)</td>
<td>3</td>
</tr>
<tr>
<td>numnode(tsquery)</td>
<td>integer</td>
<td>number of lexemes plus operators in tsquery</td>
<td>numnode('fat &amp; rat')</td>
<td>5</td>
</tr>
<tr>
<td>plainto_tsquery([])</td>
<td>tsquery</td>
<td>produce tsquery ignoring punctuation</td>
<td>plainto_tsquery('english':fat' &amp; 'rat')</td>
<td>'fat' &amp; 'rat'</td>
</tr>
<tr>
<td>phraseto_tsquery([])</td>
<td>tsquery</td>
<td>produce tsquery that searches for a phrase, ignoring punctuation</td>
<td>phraseto_tsquery('english':fat' &amp; 'rat')</td>
<td>'fat' &amp; 'rat'</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>------------------------</td>
<td>-------------</td>
<td>-------------------------------------------------------</td>
<td>-----------------------------</td>
<td>-------------------</td>
</tr>
<tr>
<td>querytree(query tsquery)</td>
<td>text</td>
<td>get indexable part of a tsquery</td>
<td>querytree('foo &amp; ! bar'::tsquery)</td>
<td>'foo'</td>
</tr>
<tr>
<td>setweight(vector tsvector, weight &quot;char&quot;)</td>
<td>tsvector</td>
<td>assign weight to each element of vector</td>
<td>setweight('fat cat:3 rat:5B'::tsvector, 'A')</td>
<td>fat':3A, fat':2A,4A, rat':5A</td>
</tr>
<tr>
<td>setweight(vector tsvector, weight &quot;char&quot;, lexemes text[])</td>
<td>tsvector</td>
<td>assign weight to elements of vector that are listed in lexemes</td>
<td>setweight('fat cat:3 rat:5B'::tsvector, 'A', '{cat, rat}')</td>
<td>fat':3A, fat':2,4, rat':5A</td>
</tr>
<tr>
<td>strip(tsvector)</td>
<td>tsvector</td>
<td>remove positions and weights from tsvector</td>
<td>strip('fat:2,4 cat:3 rat:5A'::tsvector)</td>
<td>cat':3A, fat':2,4, rat':5A</td>
</tr>
<tr>
<td>to_tsquery([config regconfig, ], query text)</td>
<td>tsquery</td>
<td>normalize words and convert to tsquery</td>
<td>to_tsquery('english', 'The &amp; Fat &amp; Rats')</td>
<td>'fat','rat'</td>
</tr>
<tr>
<td>to_tsvector([config regconfig, ], document text)</td>
<td>tsvector</td>
<td>reduce document text to tsvector</td>
<td>to_tsvector('english', 'The Fat Rats')</td>
<td>'fat','rat'</td>
</tr>
<tr>
<td>ts_delete(vector tsvector, lexeme text)</td>
<td>tsvector</td>
<td>remove given lexeme from vector</td>
<td>ts_delete('fat cat:3 rat:5A'::tsvector, 'fat')</td>
<td>fat':3, cat:3, rat:5A</td>
</tr>
<tr>
<td>ts_delete(vector tsvector, lexemes text[])</td>
<td>tsvector</td>
<td>remove any occurrence of lexemes in vector</td>
<td>ts_delete('fat cat:3 rat:5A'::tsvector, ARRAY['fat','rat'])</td>
<td>fat':3, cat:3, rat:5A</td>
</tr>
<tr>
<td>ts_filter(vector tsvector, weights &quot;char&quot;[])</td>
<td>tsvector</td>
<td>select only elements with given weights from vector</td>
<td>ts_filter('fat cat:3 rat:5A'::tsvector, ARRAY['a,b'])</td>
<td>fat':3B, cat:3b, rat:5A</td>
</tr>
</tbody>
</table>
| ts_headline([config regconfig, ], document text, query tsquery [, options text ]}) | text        | display a query match                                   | ts_headline('x y <b>z</b>', 'z':tsquery) | x y <b>z</b>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ts_rank([weights float4[], ] vector tsvector, query tsquery [, normalization integer ])</code></td>
<td>float4</td>
<td>rank document for query</td>
<td><code>ts_rank(textsearch, query)</code></td>
<td>0.818</td>
</tr>
<tr>
<td><code>ts_rank_cd([weights float4[], ] vector tsvector, query tsquery [, normalization integer ])</code></td>
<td>float4</td>
<td>rank document for query using cover density</td>
<td><code>ts_rank_cd('0 0.2, 0.4, 1.0', textsearch, query)</code></td>
<td>2.01317</td>
</tr>
<tr>
<td><code>ts_rewrite(query tsquery, target tsquery, substitute tsquery)</code></td>
<td>tsquery</td>
<td>replace target with substitute within query</td>
<td>`ts_rewrite('a &amp; b':tsquery, 'a':tsquery, 'foo</td>
<td>bar':tsquery)`</td>
</tr>
<tr>
<td><code>ts_rewrite(query tsquery, select text)</code></td>
<td>tsquery</td>
<td>replace using targets and substitutes from a SELECT command</td>
<td><code>SELECT ts_rewrite('a &amp; b':tsquery, 'SELECT t,s FROM aliases')</code></td>
<td>'b' &amp; ('foo</td>
</tr>
<tr>
<td><code>tsquery_phrase(query1 tsquery, query2 tsquery)</code></td>
<td>tsquery</td>
<td>make query that searches for query1 followed by query2 (same as <code>&lt;&gt;</code> operator)</td>
<td><code>tsquery_phrase(to_tsquery('fat'), to_tsquery('cat'))</code></td>
<td>'fat'&lt;&gt; 'cat'</td>
</tr>
<tr>
<td><code>tsquery_phrase(query1 tsquery, query2 tsquery, distance integer)</code></td>
<td>tsquery</td>
<td>make query that searches for query1 followed by query2 at distance distance</td>
<td><code>tsquery_phrase(to_tsquery('fat'), to_tsquery('cat'), 10)</code></td>
<td>'fat'&lt;&gt; 'cat'</td>
</tr>
<tr>
<td><code>tsvector_to_array(tsvector)</code></td>
<td>text[]</td>
<td>convert tsvector to array of lexemes</td>
<td>`tsvector_to_array('fat</td>
<td>fant2</td>
</tr>
<tr>
<td><code>tsvector_update_trigger()</code></td>
<td>trigger</td>
<td>trigger function for automatic tsvector column update</td>
<td><code>CREATE TRIGGER ... tsvector_update_trigger(tsvector, 'pg_catalog.swedish', title, body)</code></td>
<td></td>
</tr>
</tbody>
</table>
## Chapter 9. Functions and Operators

### Function | Return Type | Description | Example | Result |
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td><code>tsvector_update_trigger_column()</code></td>
<td>trigger</td>
<td>trigger function for automatic tsvector column update</td>
<td>CREATE TRIGGER ... tsvector_update_trigger_column(tsvcol, configcol, title, body)</td>
<td></td>
</tr>
<tr>
<td><code>unnest(tsvector, OUT lexeme text, OUT positions smallint[], OUT weights text)</code></td>
<td>setof record</td>
<td>expand a tsvector to a set of rows</td>
<td>unnest('fat:2,4 cat:3 rat:5A'::tsvector)</td>
<td>(cat,{3},{D}) ...</td>
</tr>
</tbody>
</table>

**Note:** All the text search functions that accept an optional `regconfig` argument will use the configuration specified by `default_text_search_config` when that argument is omitted.

The functions in Table 9-41 are listed separately because they are not usually used in everyday text searching operations. They are helpful for development and debugging of new text search configurations.

### Table 9-41. Text Search Debugging Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ts_debug([config regconfig, ] document text, OUT alias text, OUT description text, OUT token text, OUT dictionaries regdictionary[], OUT dictionary regdictionary, OUT lexemes text[])</code></td>
<td>setof record</td>
<td>test a configuration</td>
<td><code>ts_debug('english', 'The Brightest supernovaes')</code></td>
<td>{}</td>
</tr>
<tr>
<td><code>ts_lexize(dict regdictionary, token text)</code></td>
<td>text[]</td>
<td>test a dictionary</td>
<td><code>ts_lexize('english_stem', 'stars')</code></td>
<td>{star}</td>
</tr>
<tr>
<td><code>ts_parse(parser_name text, document text, OUT tokid integer, OUT token text)</code></td>
<td>setof record</td>
<td>test a parser</td>
<td><code>ts_parse('default', 'foo - bar')</code></td>
<td>{}</td>
</tr>
</tbody>
</table>

259
9.14. XML Functions

The functions and function-like expressions described in this section operate on values of type xml. Check Section 8.13 for information about the xml type. The function-like expressions xmlparse and xmlserialize for converting to and from type xml are not repeated here. Use of most of these functions requires the installation to have been built with configure --with-libxml.

9.14.1. Producing XML Content

A set of functions and function-like expressions are available for producing XML content from SQL data. As such, they are particularly suitable for formatting query results into XML documents for processing in client applications.

9.14.1.1. xmlcomment

xmlcomment (text)

The function xmlcomment creates an XML value containing an XML comment with the specified text as content. The text cannot contain "--" or end with a "-" so that the resulting construct is a valid XML comment. If the argument is null, the result is null.

Example:
SELECT xmlcomment('hello');

xmlcomment
---------------
<!--hello-->

9.14.1.2. xmlconcat

xmconcat(xml[, ...])

The function xmlconcat concatenates a list of individual XML values to create a single value containing an XML content fragment. Null values are omitted; the result is only null if there are no nonnull arguments.

Example:

SELECT xmlconcat('<abc/>', '<bar>foo</bar>');

xmlconcat
----------------------
<abc/><bar>foo</bar>

XML declarations, if present, are combined as follows. If all argument values have the same XML version declaration, that version is used in the result, else no version is used. If all argument values have the standalone declaration value "yes", then that value is used in the result. If all argument values have a standalone declaration value and at least one is "no", then that is used in the result. Else the result will have no standalone declaration. If the result is determined to require a standalone declaration but no version declaration, a version declaration with version 1.0 will be used because XML requires an XML declaration to contain a version declaration. Encoding declarations are ignored and removed in all cases.

Example:

SELECT xmlconcat('<?xml version="1.1"?><foo/>', '<?xml version="1.1" standalone="no"?><bar/>');

xmlconcat
-----------------------------------
<?xml version="1.1"?><foo/><bar/>

9.14.1.3. xmlelement

xmlelement(name name [, xmlattributes(value [AS attname] [, ... ])] [, content, ...])

The xmlelement expression produces an XML element with the given name, attributes, and content.

Examples:

SELECT xmlelement(name foo);

xmlelement
Chapter 9. Functions and Operators

---

<xoo/>

SELECT xmlelement(name foo, xmlattributes('xyz' as bar));

xmlelement

---

<xoo bar="xyz"/>

SELECT xmlelement(name foo, xmlattributes(current_date as bar), 'cont', 'ent');

xmlelement

---

<xoo bar="2007-01-26">content</foo>

Element and attribute names that are not valid XML names are escaped by replacing the offending characters by the sequence _xHHHH_, where HHHH is the character's Unicode codepoint in hexadecimal notation. For example:

SELECT xmlelement(name "foo$bar", xmlattributes('xyz' as "a&b"));

xmlelement

---

<xoo_x0024_bar a_x0026_b="xyz"/>

An explicit attribute name need not be specified if the attribute value is a column reference, in which case the column’s name will be used as the attribute name by default. In other cases, the attribute must be given an explicit name. So this example is valid:

CREATE TABLE test (a xml, b xml);
SELECT xmlelement(name test, xmlattributes(a, b)) FROM test;

But these are not:

SELECT xmlelement(name test, xmlattributes('constant'), a, b) FROM test;
SELECT xmlelement(name test, xmlattributes(func(a, b))) FROM test;

Element content, if specified, will be formatted according to its data type. If the content is itself of type xml, complex XML documents can be constructed. For example:

SELECT xmlelement(name foo, xmlattributes('xyz' as bar),
    xmlelement(name abc),
    xmlcomment('test'),
    xmlelement(name xyz));

xmlelement

---

<xoo bar="xyz"><abc/><!--test--><xyz/></foo>

Content of other types will be formatted into valid XML character data. This means in particular that the characters <, >, and & will be converted to entities. Binary data (data type bytea) will be represented in base64 or hex encoding, depending on the setting of the configuration parameter
xmlbinary. The particular behavior for individual data types is expected to evolve in order to align the SQL and PostgreSQL data types with the XML Schema specification, at which point a more precise description will appear.

### 9.14.1.4. xmlforest

```sql
xmlforest(content [AS name] [, ...])
```

The `xmlforest` expression produces an XML forest (sequence) of elements using the given names and content.

**Examples:**

```sql
SELECT xmlforest('abc' AS foo, 123 AS bar);
```

```xml
<foo>abc</foo><bar>123</bar>
```

```sql
SELECT xmlforest(table_name, column_name)
FROM information_schema.columns
WHERE table_schema = 'pg_catalog';
```

```xml
<table_name>pg_authid</table_name><column_name>rolname</column_name>
<table_name>pg_authid</table_name><column_name>rolsuper</column_name>
...
```

As seen in the second example, the element name can be omitted if the content value is a column reference, in which case the column name is used by default. Otherwise, a name must be specified.

Element names that are not valid XML names are escaped as shown for `xmlelement` above. Similarly, content data is escaped to make valid XML content, unless it is already of type `xml`.

Note that XML forests are not valid XML documents if they consist of more than one element, so it might be useful to wrap `xmlforest` expressions in `xmlelement`.

### 9.14.1.5. xmlpi

```sql
xmlpi(name target [, content])
```

The `xmlpi` expression creates an XML processing instruction. The content, if present, must not contain the character sequence `?>`.

**Example:**

```sql
SELECT xmlpi(name php, 'echo "hello world";');
```

```xml
<?php echo "hello world";?>
```
9.14.1.6.\_xmlroot

xmlroot\(xml, \text{version text | no value [, standalone yes|no|no value]}\)

The `xmlroot` expression alters the properties of the root node of an XML value. If a version is specified, it replaces the value in the root node’s version declaration; if a standalone setting is specified, it replaces the value in the root node’s standalone declaration.

```sql
SELECT xmlroot(xmlparse(document '<?xml version="1.1"?><content>abc</content>'),
    version '1.0', standalone yes);
```

```
<?xml version="1.0" standalone="yes"?><content>abc</content>
```

9.14.1.7.\_xmlagg

xmlagg\(xml\)

The function `xmlagg` is, unlike the other functions described here, an aggregate function. It concatenates the input values to the aggregate function call, much like `xmlconcat` does, except that concatenation occurs across rows rather than across expressions in a single row. See Section 9.20 for additional information about aggregate functions.

Example:

```sql
CREATE TABLE test (y int, x xml);
INSERT INTO test VALUES (1, '<foo>abc</foo>');
INSERT INTO test VALUES (2, '<bar/>');
SELECT xmlagg(x) FROM test;
```

```
<foo>abc</foo><bar/>
```

To determine the order of the concatenation, an `ORDER BY` clause may be added to the aggregate call as described in Section 4.2.7. For example:

```sql
SELECT xmlagg(x ORDER BY y DESC) FROM test;
```

```
<bar/><foo>abc</foo>
```

The following non-standard approach used to be recommended in previous versions, and may still be useful in specific cases:

```sql
SELECT xmlagg(x) FROM (SELECT * FROM test ORDER BY y DESC) AS tab;
```

```
<bar/><foo>abc</foo>
```
9.14.2. XML Predicates

The expressions described in this section check properties of xml values.

9.14.2.1. IS DOCUMENT

`IS DOCUMENT` returns true if the argument XML value is a proper XML document, false if it is not (that is, it is a content fragment), or null if the argument is null. See Section 8.13 about the difference between documents and content fragments.

9.14.2.2. IS NOT DOCUMENT

`IS NOT DOCUMENT` returns false if the argument XML value is a proper XML document, true if it is not (that is, it is a content fragment), or null if the argument is null.

9.14.2.3. XMLEXISTS

`XMLEXISTS(text PASSING [BY REF] xml [BY REF])` returns true if the XPath expression in the first argument returns any nodes, and false otherwise. (If either argument is null, the result is null.)

Example:

```sql
SELECT xmlexists('//town[text() = "Toronto"]' PASSING BY REF '<towns><town>Toronto</town><town>Ottawa</town></towns>');
```

```
<table>
<thead>
<tr>
<th>xmlexists</th>
</tr>
</thead>
<tbody>
<tr>
<td>t</td>
</tr>
</tbody>
</table>
```

The BY REF clauses have no effect in PostgreSQL, but are allowed for SQL conformance and compatibility with other implementations. Per SQL standard, the first BY REF is required, the second is optional. Also note that the SQL standard specifies the xmlexists construct to take an XQuery expression as first argument, but PostgreSQL currently only supports XPath, which is a subset of XQuery.

9.14.2.4. xml_is_well_formed

`xml_is_well_formed(text)`
`xml_is_well_formed_document(text)`
`xml_is_well_formed_content(text)`
These functions check whether a text string is well-formed XML, returning a Boolean result. `xml_is_well_formed_document` checks for a well-formed document, while `xml_is_well_formed_content` checks for well-formed content. `xml_is_well_formed` does the former if the `xmloption` configuration parameter is set to DOCUMENT, or the latter if it is set to CONTENT. This means that `xml_is_well_formed` is useful for seeing whether a simple cast to type `xml` will succeed, whereas the other two functions are useful for seeing whether the corresponding variants of `XMLPARSE` will succeed.

Examples:

```sql
SET xmloption TO DOCUMENT;
SELECT xml_is_well_formed(' <>');

xml_is_well_formed
-----------------------------
 f   (1 row)

SELECT xml_is_well_formed('<abc/>');

xml_is_well_formed
-----------------------------
 t   (1 row)

SET xmloption TO CONTENT;
SELECT xml_is_well_formed('abc');

xml_is_well_formed
-----------------------------
 t   (1 row)

SELECT xml_is_well_formed_document('<pg:foo xmlns:pg="http://postgresql.org/stuff">bar</pg:foo>');

xml_is_well_formed_document
-----------------------------------------------
 t   (1 row)

SELECT xml_is_well_formed_document('<pg:foo xmlns:pg="http://postgresql.org/stuff">bar</my:foo>');

xml_is_well_formed_document
-----------------------------------------------
 f   (1 row)
```

The last example shows that the checks include whether namespaces are correctly matched.

### 9.14.3. Processing XML

To process values of data type `xml`, PostgreSQL offers the functions `xpath` and `xpath_exists`, which evaluate XPath 1.0 expressions.

```
xpath(xpath, xml [, nsarray])
```

The function `xpath` evaluates the XPath expression `xpath` (a text value) against the XML value `xml`. It returns an array of XML values corresponding to the node set produced by the XPath expression. If the XPath expression returns a scalar value rather than a node set, a single-element array is returned.
The second argument must be a well formed XML document. In particular, it must have a single root node element.

The optional third argument of the function is an array of namespace mappings. This array should be a two-dimensional text array with the length of the second axis being equal to 2 (i.e., it should be an array of arrays, each of which consists of exactly 2 elements). The first element of each array entry is the namespace name (alias), the second the namespace URI. It is not required that aliases provided in this array be the same as those being used in the XML document itself (in other words, both in the XML document and in the xpath function context, aliases are local).

Example:

```sql
    ARRAY[ARRAY['my', 'http://example.com']]);
```

```
xpath
--------
{test}
(1 row)
```

To deal with default (anonymous) namespaces, do something like this:

```sql
SELECT xpath('//mydefns:b/text()', '<a xmlns="http://example.com"><b>test</b></a>',
    ARRAY[ARRAY['mydefns', 'http://example.com']]);
```

```
xpath
--------
{test}
(1 row)
```

`xpath_exists`(`xpath`, `xml` [, `nsarray`])

The function `xpath_exists` is a specialized form of the `xpath` function. Instead of returning the individual XML values that satisfy the XPath, this function returns a Boolean indicating whether the query was satisfied or not. This function is equivalent to the standard XMLEXISTS predicate, except that it also offers support for a namespace mapping argument.

Example:

```sql
    ARRAY[ARRAY['my', 'http://example.com']]);
```

```
xpath_exists
-------------
t
(1 row)
```
9.14.4. Mapping Tables to XML

The following functions map the contents of relational tables to XML values. They can be thought of as XML export functionality:

- `table_to_xml(tbl regclass, nulls boolean, tableforest boolean, targetns text)`
- `query_to_xml(query text, nulls boolean, tableforest boolean, targetns text)`
- `cursor_to_xml(cursor refcursor, count int, nulls boolean, tableforest boolean, targetns text)`

The return type of each function is `xml`.

`table_to_xml` maps the content of the named table, passed as parameter `tbl`. The `regclass` type accepts strings identifying tables using the usual notation, including optional schema qualifications and double quotes. `query_to_xml` executes the query whose text is passed as parameter `query` and maps the result set. `cursor_to_xml` fetches the indicated number of rows from the cursor specified by the parameter `cursor`. This variant is recommended if large tables have to be mapped, because the result value is built up in memory by each function.

If `tableforest` is false, then the resulting XML document looks like this:

```xml
<tablename>
  <row>
    <columnname1>data</columnname1>
    <columnname2>data</columnname2>
  </row>
  ...
</tablename>
```

If `tableforest` is true, the result is an XML content fragment that looks like this:

```xml
<tablename>
  <columnname1>data</columnname1>
  <columnname2>data</columnname2>
</tablename>
<tablename>
  ...
</tablename>
```

If no table name is available, that is, when mapping a query or a cursor, the string `table` is used in the first format, `row` in the second format.

The choice between these formats is up to the user. The first format is a proper XML document, which will be important in many applications. The second format tends to be more useful in the `cursor_to_xml` function if the result values are to be reassembled into one document later on. The functions for producing XML content discussed above, in particular `xmlelement`, can be used to alter the results to taste.

The data values are mapped in the same way as described for the function `xmlelement` above.
The parameter `nulls` determines whether null values should be included in the output. If true, null values in columns are represented as:

\[
\text{<columnname xsi:nil="true"/>}
\]

where `xsi` is the XML namespace prefix for XML Schema Instance. An appropriate namespace declaration will be added to the result value. If false, columns containing null values are simply omitted from the output.

The parameter `targetns` specifies the desired XML namespace of the result. If no particular namespace is wanted, an empty string should be passed.

The following functions return XML Schema documents describing the mappings performed by the corresponding functions above:

- `table_to_xmlschema(tbl regclass, nulls boolean, tableforest boolean, targetns text)`
- `query_to_xmlschema(query text, nulls boolean, tableforest boolean, targetns text)`
- `cursor_to_xmlschema(cursor refcursor, nulls boolean, tableforest boolean, targetns text)`

It is essential that the same parameters are passed in order to obtain matching XML data mappings and XML Schema documents.

The following functions produce XML data mappings and the corresponding XML Schema in one document (or forest), linked together. They can be useful where self-contained and self-describing results are wanted:

- `table_to_xml_and_xmlschema(tbl regclass, nulls boolean, tableforest boolean, targetns text)`
- `query_to_xml_and_xmlschema(query text, nulls boolean, tableforest boolean, targetns text)`

In addition, the following functions are available to produce analogous mappings of entire schemas or the entire current database:

- `schema_to_xml(schema name, nulls boolean, tableforest boolean, targetns text)`
- `schema_to_xmlschema(schema name, nulls boolean, tableforest boolean, targetns text)`
- `schema_to_xml_and_xmlschema(schema name, nulls boolean, tableforest boolean, targetns text)`
- `database_to_xml(nulls boolean, tableforest boolean, targetns text)`
- `database_to_xmlschema(nulls boolean, tableforest boolean, targetns text)`
- `database_to_xml_and_xmlschema(nulls boolean, tableforest boolean, targetns text)`

Note that these potentially produce a lot of data, which needs to be built up in memory. When requesting content mappings of large schemas or databases, it might be worthwhile to consider mapping the tables separately instead, possibly even through a cursor.

The result of a schema content mapping looks like this:

\[
\text{<schemaname>}
\text{table1-mapping}
\text{table2-mapping}
\ldots
\text{</schemaname>}
\]

where the format of a table mapping depends on the `tableforest` parameter as explained above.
The result of a database content mapping looks like this:

```xml
<dbname>
  <schema1name>...
  </schema1name>
  <schema2name>...
  </schema2name>
  ...
</dbname>
```

where the schema mapping is as above.

As an example of using the output produced by these functions, Figure 9-1 shows an XSLT stylesheet that converts the output of `table_to_xml_and_xmlschema` to an HTML document containing a tabular rendition of the table data. In a similar manner, the results from these functions can be converted into other XML-based formats.

**Figure 9-1. XSLT Stylesheet for Converting SQL/XML Output to HTML**

```xml
<?xml version="1.0"?>
<xsl:stylesheet version="1.0"
  xmlns:xsl="http://www.w3.org/1999/XSL/Transform"
  xmlns:xsd="http://www.w3.org/2001/XMLSchema"
  xmlns="http://www.w3.org/1999/xhtml">
  <xsl:output method="xml"
    doctype-system="http://www.w3.org/TR/xhtml1/DTD/xhtml1-strict.dtd"
    doctype-public="-//W3C/DTD XHTML 1.0 Strict//EN"
    indent="yes"/>

  <xsl:template match="/">
    <xsl:variable name="schema" select="//xsd:schema"/>
    <xsl:variable name="tabletypename" select="$schema/xsd:element[@name=name(current())]/@type"/>
    <xsl:variable name="rowtypename" select="$schema/xsd:complexType[@name=$tabletypename]/xsd:sequence/xsd:element[@name='row']/@type"/>

    <html>
      <head>
        <title><xsl:value-of select="name(current())"></title>
      </head>
      <body>
        <table>
          <tr>
            <xsl:for-each select="$schema/xsd:complexType[@name=$tabletypename]/xsd:sequence/xsd:element[@name='row']/@name">
              <th><xsl:value-of select="." /></th>
            </xsl:for-each>
          </tr>

          <xsl:for-each select="row">
            <tr>
              <xsl:for-each select="$schema/xsd:complexType[@name=$rowtypename]/xsd:sequence/xsd:element[@name='row']/@name">
                <td><xsl:value-of select="." /></td>
              </xsl:for-each>
            </tr>
          </xsl:for-each>
        </table>
      </body>
    </html>
  </xsl:template>
</xsl:stylesheet>
```
9.15. JSON Functions and Operators

Table 9-42 shows the operators that are available for use with the two JSON data types (see Section 8.14).

<table>
<thead>
<tr>
<th>Operator</th>
<th>RightOperandType</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>&gt;</code></td>
<td>int</td>
<td>Get JSON array element (indexed from zero, negative integers count from the end)</td>
<td><code>{&quot;a&quot;:&quot;foo&quot;}, </code>{b&quot;:&quot;baz}::json&gt;`</td>
<td><code>{c&quot;: &quot;baz\}::json&gt;</code></td>
</tr>
<tr>
<td><code>&gt;</code></td>
<td>text</td>
<td>Get JSON object field by key</td>
<td><code>{a&quot;:{b&quot;:&quot;foo\}}::json&gt;</code></td>
<td><code>{b&quot;:&quot;foo\}::json&gt;</code></td>
</tr>
<tr>
<td><code>&gt;&gt;&gt;</code></td>
<td>int</td>
<td>Get JSON array element as text</td>
<td>`{1,2,3}::json&gt;&gt;2</td>
<td></td>
</tr>
<tr>
<td><code>&gt;&gt;&gt;</code></td>
<td>text</td>
<td>Get JSON object field as text</td>
<td>`{a&quot;:1,&quot;b&quot;:2}&quot;::json&gt;&gt;&gt;'b'</td>
<td></td>
</tr>
<tr>
<td><code>#</code></td>
<td>text[]</td>
<td>Get JSON object at specified path</td>
<td>`{a&quot;:{c&quot;: &quot;foo})::json#&gt;'{a,b}'</td>
<td></td>
</tr>
<tr>
<td><code>&gt;&gt;&gt;</code></td>
<td>text[]</td>
<td>Get JSON object at specified path as text</td>
<td>`{&quot;a&quot;:{1,2,3&quot;,&quot;b&quot;:{4,5,6}::json#&gt;'{a,2}'</td>
<td></td>
</tr>
</tbody>
</table>

**Note:** There are parallel variants of these operators for both the `json` and `jsonb` types. The field/element/path extraction operators return the same type as their left-hand input (either `json` or `jsonb`), except for those specified as returning `text`, which coerce the value to text. The field/element/path extraction operators return NULL, rather than failing, if the JSON input does not have the right structure to match the request; for example if no such element exists. The field/element/path extraction operators that accept integer JSON array subscripts all support neg-
ative subscripting from the end of arrays.

The standard comparison operators shown in Table 9-1 are available for jsonb, but not for json. They follow the ordering rules for B-tree operations outlined at Section 8.14.4.

Some further operators also exist only for jsonb, as shown in Table 9-43. Many of these operators can be indexed by jsonb operator classes. For a full description of jsonb containment and existence semantics, see Section 8.14.3. Section 8.14.4 describes how these operators can be used to effectively index jsonb data.

Table 9-43. Additional jsonb Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Right Operand Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>@&gt;</td>
<td>jsonb</td>
<td>Does the left JSON value contain the right JSON path/value entries at the top level?</td>
<td>'{&quot;a&quot;:1, &quot;b&quot;:2}':jsonb @&gt; '{&quot;b&quot;:2}':jsonb</td>
</tr>
<tr>
<td>&lt;@</td>
<td>jsonb</td>
<td>Are the left JSON path/value entries contained at the top level within the right JSON value?</td>
<td>'{&quot;b&quot;:2}':jsonb &lt;@ '{&quot;a&quot;:1, &quot;b&quot;:2}':jsonb</td>
</tr>
<tr>
<td>?</td>
<td>text</td>
<td>Does the string exist as a top-level key within the JSON value?</td>
<td>'{&quot;a&quot;:1, &quot;b&quot;:2}'::jsonb ? 'b'</td>
</tr>
<tr>
<td>?_</td>
<td>text[]</td>
<td>Do any of these array strings exist as top-level keys?</td>
<td>'{&quot;a&quot;:1, &quot;b&quot;:2, &quot;c&quot;:3}':jsonb ?_ array['b', 'c']</td>
</tr>
<tr>
<td>?&amp;</td>
<td>text[]</td>
<td>Do all of these array strings exist as top-level keys?</td>
<td>'{&quot;a&quot;, &quot;b&quot;}':jsonb ?&amp; array['a', 'b']</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>jsonb</td>
</tr>
<tr>
<td>-</td>
<td>text</td>
<td>Delete key/value pair or string element from left operand. Key/value pairs are matched based on their key value.</td>
<td>'{&quot;a&quot;:1, &quot;b&quot;}':jsonb - 'a'</td>
</tr>
<tr>
<td>-</td>
<td>integer</td>
<td>Delete the array element with specified index (Negative integers count from the end). Throws an error if top level container is not an array.</td>
<td>'{&quot;a&quot;:1, &quot;b&quot;}':jsonb - 1</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Right Operand Type</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
</table>
| #-       | text[]            | Delete the field or element with specified path (for JSON arrays, negative integers count from the end) | ‘{"a", {"b":1}}’::jsonb 
#- ‘(1,b)’ |

**Note:** The || operator concatenates the elements at the top level of each of its operands. It does not operate recursively. For example, if both operands are objects with a common key field name, the value of the field in the result will just be the value from the right hand operand.

Table 9-44 shows the functions that are available for creating json and jsonb values. (There are no equivalent functions for jsonb, of the row_to_json and array_to_json functions. However, the to_jsonb function supplies much the same functionality as these functions would.)

Table 9-44. JSON Creation Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>to_json(anyelement)</td>
<td>Returns the value as json or jsonb. Arrays and composites are converted (recursively) to arrays and objects; otherwise, if there is a cast from the type to json, the cast function will be used to perform the conversion; otherwise, a scalar value is produced. For any scalar type other than a number, a Boolean, or a null value, the text representation will be used, in such a fashion that it is a valid json or jsonb value.</td>
<td>to_json('Fred said &quot;Hi.&quot;'::text)</td>
<td>&quot;Fred said &quot;Hi.&quot;&quot;</td>
</tr>
<tr>
<td>array_to_json(anyarray [, pretty_bool])</td>
<td>Returns the array as a JSON array. A PostgreSQL multidimensional array becomes a JSON array of arrays. Line feeds will be added between dimension-1 elements if pretty_bool is true.</td>
<td>array_to_json('{{1,5}{</td>
<td>99,100}}'::int[])</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>row_to_json(record [, pretty_bool])</td>
<td>Returns the row as a JSON object. Line feeds will be added between level-1 elements if pretty_bool is true.</td>
<td>row_to_json(row(1, &quot;foo&quot;))</td>
<td>&quot;f1&quot;:1,&quot;f2&quot;:&quot;foo&quot;</td>
</tr>
<tr>
<td>json_build_array(VARIADIC &quot;any&quot;)</td>
<td>Builds a possibly-heterogeneously-typed JSON array out of a variadic argument list.</td>
<td>json_build_array(1, 2, 3, 4, 5)</td>
<td>[1, 2, &quot;3&quot;, 4, 5]</td>
</tr>
<tr>
<td>json_build_object(VARIADIC &quot;any&quot;)</td>
<td>Builds a JSON object out of a variadic argument list. By convention, the argument list consists of alternating keys and values.</td>
<td>json_build_object(&quot;foo&quot;, 1, &quot;bar&quot;, 2)</td>
<td>{&quot;foo&quot;: 1, &quot;bar&quot;: 2}</td>
</tr>
<tr>
<td>json_object(text[])</td>
<td>Builds a JSON object out of a text array. The array must have either exactly one dimension with an even number of members, in which case they are taken as alternating key/value pairs, or two dimensions such that each inner array has exactly two elements, which are taken as a key/value pair.</td>
<td>json_object('{a, 1, b, &quot;def&quot;, c, 3.5}')</td>
<td>{&quot;a&quot;: &quot;1&quot;, &quot;b&quot;: &quot;def&quot;, &quot;c&quot;: &quot;3.5&quot;}</td>
</tr>
<tr>
<td>json_object(keys text[], values text[])</td>
<td>This form of json_object takes keys and values pairwise from two separate arrays. In all other respects it is identical to the one-argument form.</td>
<td>json_object('a, b', '(1,2)')</td>
<td>{&quot;a&quot;: &quot;1&quot;, &quot;b&quot;: &quot;2&quot;}</td>
</tr>
</tbody>
</table>

**Note:** array_to_json and row_to_json have the same behavior as to_json except for offering a pretty-printing option. The behavior described for to_json likewise applies to each individual value converted by the other JSON creation functions.

**Note:** The hstore extension has a cast from hstore to json, so that hstore values converted via
Chapter 9. Functions and Operators
the JSON creation functions will be represented as JSON objects, not as primitive string values.

Table 9-45 shows the functions that are available for processing json and jsonb values.
Table 9-45. JSON Processing Functions
Function

Return Type

Description

Example

json_array_length(json)
int
jsonb_array_length(jsonb)

Returns the
number of
elements in the
outermost JSON
array.

json_array_length(’[1,2,3,{"f1":1,"f2":[5,6
5

json_each(json)setof
jsonb_each(jsonb)
text,
json
setof
text,
jsonb

Expands the
select * from
key | value
outermost JSON json_each(’{"a":"foo",
-----+------object into a set of "b":"bar"}’)
a
| "foo"
key/value pairs.
b
| "bar"

key
value
key
value

Example Result

json_each_text(json)
setof key
jsonb_each_text(jsonb)
text, value
text

Expands the
select * from
key | value
outermost JSON json_each_text(’{"a":"foo",
-----+------object into a set of "b":"bar"}’)
a
| foo
key/value pairs.
b
| bar
The returned
values will be of
type text.

json_extract_path(from_json
json
json,
jsonb
VARIADIC
path_elems
text[])
jsonb_extract_path(from_json
jsonb,
VARIADIC
path_elems
text[])

Returns JSON
value pointed to
by path_elems
(equivalent to #>
operator).

json_extract_path(’{"f2":{"f3":1},"f4":{"f5
{"f5":99,"f6":"foo"}

275


Chapter 9. Functions and Operators
Function

Return Type

Description

Example

Example Result

json_extract_path_text(from_json
text
Returns JSON
json,
value pointed to
VARIADIC
by path_elems
path_elems
as text
text[])
(equivalent to
jsonb_extract_path_text(from_json
#>> operator).
jsonb,
VARIADIC
path_elems
text[])

json_extract_path_text(’{"f2":{"f3":1},"f4"
foo
’f6’)

json_object_keys(json)
setof text
jsonb_object_keys(jsonb)

Returns set of
keys in the
outermost JSON
object.

json_object_keys(’{"f1":"abc","f2":{"f3":"a
json_object_keys
"f4":"b"}}’)
-----------------f1 f2

json_populate_record(base
anyelement
anyelement,
from_json
json)
jsonb_populate_record(base
anyelement,
from_json
jsonb)

Expands the
object in
from_json to a
row whose
columns match
the record type
defined by base
(see note below).

select * from
a | b
json_populate_record(null::myrowtype,
---+--’{"a":1,"b":2}’)
1 | 2

json_populate_recordset(base
setof
anyelement,
anyelement
from_json
json)
jsonb_populate_recordset(base
anyelement,
from_json
jsonb)

Expands the
select * from
a | b
outermost array of json_populate_recordset(null::myrowtype,
---+--objects in
’[{"a":1,"b":2},{"a":3,"b":4}]’)
1 | 2 3 | 4
from_json to a
set of rows whose
columns match
the record type
defined by base
(see note below).

json_array_elements(json)
setof json
jsonb_array_elements(jsonb)
setof jsonb

Expands a JSON
array to a set of
JSON values.

json_array_elements_text(json)
setof text
Expands a JSON
jsonb_array_elements_text(jsonb)
array to a set of
text values.

select * from
value
json_array_elements(’[1,true,
----------[2,false]]’)
1 true
[2,false]
select * from
value
json_array_elements_text(’["foo",
----------"bar"]’)
foo bar

276


Chapter 9. Functions and Operators
Function

Return Type

Description

Example

Example Result

json_typeof(json)
text
jsonb_typeof(jsonb)

Returns the type json_typeof(’-123.4’)
number
of the outermost
JSON value as a
text string.
Possible types are
object, array,
string, number,
boolean, and
null.

json_to_record(json)
record
jsonb_to_record(jsonb)

Builds an
arbitrary record
from a JSON
object (see note
below). As with
all functions
returning record,
the caller must
explicitly define
the structure of
the record with an
AS clause.

select * from
a |
b
| d
json_to_record(’{"a":1,"b":[1,2,3],"c":"bar
---+---------+--as x(a int, b 1 | [1,2,3] |
text, d text)

json_to_recordset(json)
setof record
jsonb_to_recordset(jsonb)

Builds an
arbitrary set of
records from a
JSON array of
objects (see note
below). As with
all functions
returning record,
the caller must
explicitly define
the structure of
the record with an
AS clause.

select * from
a | b
json_to_recordset(’[{"a":1,"b":"foo"},{"a":
---+----as x(a int, b 1 | foo 2 |
text);

json_strip_nulls(from_json
json
json)
jsonb
jsonb_strip_nulls(from_json
jsonb)

Returns

json_strip_nulls(’[{"f1":1,"f2":null},2,nul
[{"f1":1},2,null,3]

from_json with

all object fields
that have null
values omitted.
Other null values
are untouched.

277


<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Example Result</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>jsonb_set(target jsonb, path text[], new_value jsonb [, create_missing boolean])</code></td>
<td><code>jsonb</code></td>
<td>Returns target with the section designated by <code>path</code> replaced by <code>new_value</code>, or with <code>new_value</code> added if <code>create_missing</code> is true (default is true) and the item designated by <code>path</code> does not exist. As with the path oriented operators, negative integers that appear in <code>path</code> count from the end of JSON arrays.</td>
<td><code>jsonb_set('&quot;\{&quot;f1&quot;:1, f2&quot;:null, f3&quot;:[2,3,4]\}&quot;', '{0,f1}','[2,3,4]',false)</code></td>
<td><code>\{&quot;f1&quot;: [2,3,4], f2&quot;:null, f3&quot;: [2,3,4]\}</code></td>
</tr>
<tr>
<td><code>jsonb_insert(target jsonb, path text[], new_value jsonb [, insert_after boolean])</code></td>
<td><code>jsonb</code></td>
<td>Returns target with <code>new_value</code> inserted. If target section designated by <code>path</code> is in a JSONB array, <code>new_value</code> will be inserted before target or after if <code>insert_after</code> is true (default is false). If target section designated by <code>path</code> is in JSONB object, <code>new_value</code> will be inserted only if target does not exist. As with the path oriented operators, negative integers that appear in <code>path</code> count from the end of JSON arrays.</td>
<td><code>jsonb_insert('&quot;\{&quot;a&quot;:[0,1,2]\}&quot;', '{a, 1}','&quot;new_value&quot;')</code></td>
<td><code>{&quot;a&quot;: [0,1,2], &quot;new_value&quot;: 1, 2}</code></td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators
Function

Return Type

jsonb_pretty(from_json
text
jsonb)

Description

Example

Returns

jsonb_pretty(’[{"f1":1,"f2":null},2,null,3]
[
{
"f1": 1,
"f2": null
},
2,
null,
3 ]

from_json as

indented JSON
text.

Example Result

Note: Many of these functions and operators will convert Unicode escapes in JSON strings to the
appropriate single character. This is a non-issue if the input is type jsonb, because the conversion
was already done; but for json input, this may result in throwing an error, as noted in Section 8.14.

Note:

While
the
examples
for
the
functions
json_populate_record,
json_populate_recordset, json_to_record and json_to_recordset use constants, the
typical use would be to reference a table in the FROM clause and use one of its json or jsonb
columns as an argument to the function. Extracted key values can then be referenced in other
parts of the query, like WHERE clauses and target lists. Extracting multiple values in this way can
improve performance over extracting them separately with per-key operators.
JSON keys are matched to identical column names in the target row type. JSON type coercion for
these functions is “best effort” and may not result in desired values for some types. JSON fields
that do not appear in the target row type will be omitted from the output, and target columns that
do not match any JSON field will simply be NULL.

Note: All the items of the path parameter of jsonb_set as well as jsonb_insert except the last
item must be present in the target. If create_missing is false, all items of the path parameter
of jsonb_set must be present. If these conditions are not met the target is returned unchanged.
If the last path item is an object key, it will be created if it is absent and given the new value. If
the last path item is an array index, if it is positive the item to set is found by counting from the
left, and if negative by counting from the right - -1 designates the rightmost element, and so on.
If the item is out of the range -array_length .. array_length -1, and create_missing is true, the new
value is added at the beginning of the array if the item is negative, and at the end of the array if it
is positive.

Note: The json_typeof function’s null return value should not be confused with a SQL NULL.
While calling json_typeof(’null’::json) will return null, calling json_typeof(NULL::json)
will return a SQL NULL.

Note: If the argument to json_strip_nulls contains duplicate field names in any object, the
result could be semantically somewhat different, depending on the order in which they occur.
This is not an issue for jsonb_strip_nulls since jsonb values never have duplicate object field
names.

279


See also Section 9.20 for the aggregate function `json_agg` which aggregates record values as JSON, and the aggregate function `json_object_agg` which aggregates pairs of values into a JSON object, and their `jsonb` equivalents, `jsonb_agg` and `jsonb_object_agg`.

### 9.16. Sequence Manipulation Functions

This section describes functions for operating on sequence objects, also called sequence generators or just sequences. Sequence objects are special single-row tables created with CREATE SEQUENCE. Sequence objects are commonly used to generate unique identifiers for rows of a table. The sequence functions, listed in Table 9-46, provide simple, multiuser-safe methods for obtaining successive sequence values from sequence objects.

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>currval(regclass)</code></td>
<td>bigint</td>
<td>Return value most recently obtained with <code>nextval</code> for specified sequence</td>
</tr>
<tr>
<td><code>lastval()</code></td>
<td>bigint</td>
<td>Return value most recently obtained with <code>nextval</code> for any sequence</td>
</tr>
<tr>
<td><code>nextval(regclass)</code></td>
<td>bigint</td>
<td>Advance sequence and return new value</td>
</tr>
<tr>
<td><code>setval(regclass, bigint)</code></td>
<td>bigint</td>
<td>Set sequence’s current value</td>
</tr>
<tr>
<td><code>setval(regclass, bigint, boolean)</code></td>
<td>bigint</td>
<td>Set sequence’s current value and is_called flag</td>
</tr>
</tbody>
</table>

The sequence to be operated on by a sequence function is specified by a `regclass` argument, which is simply the OID of the sequence in the `pg_class` system catalog. You do not have to look up the OID by hand, however, since the `regclass` data type’s input converter will do the work for you. Just write the sequence name enclosed in single quotes so that it looks like a literal constant. For compatibility with the handling of ordinary SQL names, the string will be converted to lower case unless it contains double quotes around the sequence name. Thus:

```sql
nextval('foo')  operates on sequence foo
nextval('FOO')  operates on sequence foo
nextval("Foo") operates on sequence Foo
```

The sequence name can be schema-qualified if necessary:

```sql
nextval('myschema.foo')  operates on myschema.foo
nextval("myschema".foo)  same as above
nextval('foo')            searches search path for foo
```

See Section 8.18 for more information about `regclass`.

**Note:** Before PostgreSQL 8.1, the arguments of the sequence functions were of type `text`, not `regclass`, and the above-described conversion from a text string to an OID value would happen at run time during each call. For backward compatibility, this facility still exists, but internally it is now handled as an implicit coercion from `text` to `regclass` before the function is invoked.
Chapter 9. Functions and Operators

When you write the argument of a sequence function as an unadorned literal string, it becomes a constant of type `regclass`. Since this is really just an OID, it will track the originally identified sequence despite later renaming, schema reassignment, etc. This "early binding" behavior is usually desirable for sequence references in column defaults and views. But sometimes you might want "late binding" where the sequence reference is resolved at run time. To get late-binding behavior, force the constant to be stored as a `text` constant instead of `regclass`:

```
nextval('foo'::text)
```

foo is looked up at runtime

Note that late binding was the only behavior supported in PostgreSQL releases before 8.1, so you might need to do this to preserve the semantics of old applications.

Of course, the argument of a sequence function can be an expression as well as a constant. If it is a text expression then the implicit coercion will result in a run-time lookup.

The available sequence functions are:

**nextval**

Advance the sequence object to its next value and return that value. This is done atomically: even if multiple sessions execute `nextval` concurrently, each will safely receive a distinct sequence value.

If a sequence object has been created with default parameters, successive `nextval` calls will return successive values beginning with 1. Other behaviors can be obtained by using special parameters in the `CREATE SEQUENCE` command; see its command reference page for more information.

**Important:** To avoid blocking concurrent transactions that obtain numbers from the same sequence, a `nextval` operation is never rolled back; that is, once a value has been fetched it is considered used and will not be returned again. This is true even if the surrounding transaction later aborts, or if the calling query ends up not using the value. For example an `INSERT` with an `ON CONFLICT` clause will compute the to-be-inserted tuple, including doing any required `nextval` calls, before detecting any conflict that would cause it to follow the `ON CONFLICT` rule instead. Such cases will leave unused "holes" in the sequence of assigned values. Thus, PostgreSQL sequence objects cannot be used to obtain "gapless" sequences.

**currval**

Return the value most recently obtained by `nextval` for this sequence in the current session. (An error is reported if `nextval` has never been called for this sequence in this session.) Because this is returning a session-local value, it gives a predictable answer whether or not other sessions have executed `nextval` since the current session did.

**lastval**

Return the value most recently returned by `nextval` in the current session. This function is identical to `currval`, except that instead of taking the sequence name as an argument it refers to whichever sequence `nextval` was most recently applied to in the current session. It is an error to call `lastval` if `nextval` has not yet been called in the current session.

**setval**

Reset the sequence object’s counter value. The two-parameter form sets the sequence’s `last_value` field to the specified value and sets its `is_called` field to `true`, meaning that the next `nextval` will advance the sequence before returning a value. The value reported by
currval is also set to the specified value. In the three-parameter form, is_called can be set to either true or false. true has the same effect as the two-parameter form. If it is set to false, the next nextval will return exactly the specified value, and sequence advancement commences with the following nextval. Furthermore, the value reported by currval is not changed in this case. For example,

```
SELECT setval('foo', 42);           Next nextval will return 43
SELECT setval('foo', 42, true);     Same as above
SELECT setval('foo', 42, false);    Next nextval will return 42
```

The result returned by setval is just the value of its second argument.

**Important:** Because sequences are non-transactional, changes made by setval are not undone if the transaction rolls back.

---

### 9.17. Conditional Expressions

This section describes the SQL-compliant conditional expressions available in PostgreSQL.

**Tip:** If your needs go beyond the capabilities of these conditional expressions, you might want to consider writing a stored procedure in a more expressive programming language.

#### 9.17.1. CASE

The SQL CASE expression is a generic conditional expression, similar to if/else statements in other programming languages:

```
CASE WHEN condition THEN result
    [WHEN ...]
    [ELSE result]
END
```

CASE clauses can be used wherever an expression is valid. Each condition is an expression that returns a boolean result. If the condition’s result is true, the value of the CASE expression is the result that follows the condition, and the remainder of the CASE expression is not processed. If the condition’s result is not true, any subsequent WHEN clauses are examined in the same manner. If no WHEN condition yields true, the value of the CASE expression is the result of the ELSE clause. If the ELSE clause is omitted and no condition is true, the result is null.

An example:

```
SELECT * FROM test;
```

```sql
a
---
1
2
3
```
Chapter 9. Functions and Operators

SELECT a,
       CASE WHEN a=1 THEN 'one'
             WHEN a=2 THEN 'two'
             ELSE 'other'
       END
FROM test;

| a | case |
|---+-----|
| 1 | one |
| 2 | two |
| 3 | other |

The data types of all the result expressions must be convertible to a single output type. See Section 10.5 for more details.

There is a "simple" form of CASE expression that is a variant of the general form above:

```
CASE expression
       WHEN value THEN result
       [WHEN ...]
       [ELSE result]
END
```

The first expression is computed, then compared to each of the value expressions in the WHEN clauses until one is found that is equal to it. If no match is found, the result of the ELSE clause (or a null value) is returned. This is similar to the switch statement in C.

The example above can be written using the simple CASE syntax:

```
SELECT a,
       CASE a WHEN 1 THEN 'one'
              WHEN 2 THEN 'two'
              ELSE 'other'
       END
FROM test;

| a | case |
|---+-----|
| 1 | one |
| 2 | two |
| 3 | other |
```

A CASE expression does not evaluate any subexpressions that are not needed to determine the result. For example, this is a possible way of avoiding a division-by-zero failure:

```
SELECT ... WHERE CASE WHEN x <> 0 THEN y/x > 1.5 ELSE false END;
```

Note: As described in Section 4.2.14, there are various situations in which subexpressions of an expression are evaluated at different times, so that the principle that "CASE evaluates only necessary subexpressions" is not ironclad. For example a constant 1/0 subexpression will usually
result in a division-by-zero failure at planning time, even if it's within a CASE arm that would never be entered at run time.

9.17.2. COALESCE

COALESCE(value [, ...])

The COALESCE function returns the first of its arguments that is not null. Null is returned only if all arguments are null. It is often used to substitute a default value for null values when data is retrieved for display, for example:

SELECT COALESCE(description, short_description, '(none)') ...

This returns description if it is not null, otherwise short_description if it is not null, otherwise (none).

Like a CASE expression, COALESCE only evaluates the arguments that are needed to determine the result; that is, arguments to the right of the first non-null argument are not evaluated. This SQL-standard function provides capabilities similar to NVL and IFNULL, which are used in some other database systems.

9.17.3. NULLIF

NULLIF(value1, value2)

The NULLIF function returns a null value if value1 equals value2; otherwise it returns value1. This can be used to perform the inverse operation of the COALESCE example given above:

SELECT NULLIF(value, '(none)') ...

In this example, if value is (none), null is returned, otherwise the value of value is returned.

9.17.4. GREATEST and LEAST

GREATEST(value [, ...])

LEAST(value [, ...])

The GREATEST and LEAST functions select the largest or smallest value from a list of any number of expressions. The expressions must all be convertible to a common data type, which will be the type of the result (see Section 10.5 for details). NULL values in the list are ignored. The result will be NULL only if all the expressions evaluate to NULL.

Note that GREATEST and LEAST are not in the SQL standard, but are a common extension. Some other databases make them return NULL if any argument is NULL, rather than only when all are NULL.
9.18. Array Functions and Operators

Table 9-47 shows the operators available for array types.

Table 9-47. Array Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example 1</th>
<th>Example 2</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>=</td>
<td>equal</td>
<td>ARRAY[1,1,2,1,3,1] = int[]</td>
<td>ARRAY[1,2,3]</td>
<td>t</td>
</tr>
<tr>
<td>&lt;&gt;</td>
<td>not equal</td>
<td>ARRAY[1,2,3] &lt;&gt; ARRAY[1,2,4]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td>&lt;</td>
<td>less than</td>
<td>ARRAY[1,2,3] &lt; ARRAY[1,2,4]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td>&gt;</td>
<td>greater than</td>
<td>ARRAY[1,4,3] &gt; ARRAY[1,2,4]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td>&lt;=</td>
<td>less than or equal</td>
<td>ARRAY[1,2,3] &lt;= ARRAY[1,2,4]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td>&gt;=</td>
<td>greater than or equal</td>
<td>ARRAY[1,4,3] &gt;= ARRAY[1,2,4]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td>@&gt;</td>
<td>contains</td>
<td>ARRAY[1,4,3] @&gt; ARRAY[3,1,3]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td>&lt;=@</td>
<td>is contained by</td>
<td>ARRAY[2,2,7] &lt;=@ ARRAY[1,7,4,2,6]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>overlap (have elements in common)</td>
<td>ARRAY[1,4,3] &amp;&amp; ARRAY[2,1]</td>
<td>t</td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>array-to-array concatenation</td>
<td>ARRAY[1,2,3]</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>array-to-array concatenation</td>
<td>ARRAY[1,2,3]</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>element-to-array concatenation</td>
<td>ARRAY[4,5,6]</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>array-to-element concatenation</td>
<td>ARRAY[4,5,6]</td>
</tr>
</tbody>
</table>

The array ordering operators (<, >, =, etc) compare the array contents element-by-element, using the default B-tree comparison function for the element data type, and sort based on the first difference. In multidimensional arrays the elements are visited in row-major order (last subscript varies most rapidly). If the contents of two arrays are equal but the dimensionality is different, the first difference in the dimensionality information determines the sort order. (This is a change from versions of PostgreSQL prior to 8.2: older versions would claim that two arrays with the same contents were equal, even if the number of dimensions or subscript ranges were different.)

The array containment operators (<@ and @>) consider one array to be contained in another one if each of its elements appears in the other one. Duplicates are not treated specially, thus ARRAY[1] and ARRAY[1,1] are each considered to contain the other.

See Section 8.15 for more details about array operator behavior. See Section 11.2 for more details about which operators support indexed operations.

Table 9-48 shows the functions available for use with array types. See Section 8.15 for more informa-
tion and examples of the use of these functions.

Table 9-48. Array Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>array_append(anyarray,</td>
<td>anyarray</td>
<td>append an element to the end of an array</td>
<td>array_append(ARRAY[1,2], 3)</td>
<td></td>
</tr>
<tr>
<td>anyelement)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>array_cat(anyarray,</td>
<td>anyarray</td>
<td>concatenate two arrays</td>
<td>array_cat(ARRAY[1,2,3], 4)</td>
<td></td>
</tr>
<tr>
<td>anyarray)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>array_ndims(anyarray)</td>
<td>int</td>
<td>returns the number of dimensions of the array</td>
<td>array_ndims(ARRAY[1,2,3], 4)</td>
<td></td>
</tr>
<tr>
<td>array_dims(anyarray)</td>
<td>text</td>
<td>returns a text representation of array's dimensions</td>
<td>array_dims(ARRAY[1,2,3], 4)</td>
<td></td>
</tr>
<tr>
<td>array_fill(anyelement,</td>
<td>anyarray</td>
<td>returns an array initialized with supplied value and dimensions,</td>
<td>array_fill(7, ARRAY[1,2,3], ARRAY[4,5,6])</td>
<td></td>
</tr>
<tr>
<td>int[], int[])</td>
<td></td>
<td>optionally with lower bounds other than 1</td>
<td></td>
<td></td>
</tr>
<tr>
<td>array_length(anyarray,</td>
<td>int</td>
<td>returns the length of the requested array dimension</td>
<td>array_length(ARRAY[1,2,3], 4)</td>
<td></td>
</tr>
<tr>
<td>int)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>array_lower(anyarray,</td>
<td>int</td>
<td>returns lower bound of the requested array dimension</td>
<td>array_lower(ARRAY[1,2,3], 1)</td>
<td></td>
</tr>
<tr>
<td>int)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>array_position(anyarray,</td>
<td>int</td>
<td>returns the subscript of the first occurrence of the second argument in</td>
<td>array_position(ARRAY[1,2,3], 1)</td>
<td></td>
</tr>
<tr>
<td>anyelement [], int])</td>
<td></td>
<td>the array, starting at the element indicated by the third argument or</td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td>at the first element (array must be one-dimensional)</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
<td>Example</td>
<td>Result</td>
</tr>
<tr>
<td>--------------------------</td>
<td>--------------</td>
<td>-----------------------------------------------------------------------------</td>
<td>--------------------------</td>
<td>-------------------------</td>
</tr>
<tr>
<td>array_positions</td>
<td>int[]</td>
<td>returns an array of subscripts of all occurrences of the second argument in the array given as first argument (array must be one-dimensional)</td>
<td><code>array_positions(ARRAY['A','A','B','A'], 'A')</code></td>
<td><code>{1,2,4}</code></td>
</tr>
<tr>
<td>array_prepend</td>
<td>anyarray</td>
<td>append an element to the beginning of an array</td>
<td><code>array_prepend(1, ARRAY[2,3])</code></td>
<td><code>{1,2,3}</code></td>
</tr>
<tr>
<td>array_remove</td>
<td>anyarray</td>
<td>remove all elements equal to the given value from the array (array must be one-dimensional)</td>
<td><code>array_remove(ARRAY[1,2,3,2], 2)</code></td>
<td><code>{1,3}</code></td>
</tr>
<tr>
<td>array_replace</td>
<td>anyarray</td>
<td>replace each array element equal to the given value with a new value</td>
<td><code>array_replace(ARRAY[1,2,5,4], 5, 3)</code></td>
<td><code>{1,2,3,4}</code></td>
</tr>
<tr>
<td>array_to_string</td>
<td>text</td>
<td>concatenates array elements using supplied delimiter and optional null string</td>
<td><code>array_to_string(ARRAY[1,2,3, NULL, 5], ',', '*')</code></td>
<td><code>1,2,3,*,5</code></td>
</tr>
<tr>
<td>array_upper</td>
<td>int</td>
<td>returns upper bound of the requested array dimension</td>
<td><code>array_upper(ARRAY[1,8,3,7], 1)</code></td>
<td>4</td>
</tr>
<tr>
<td>cardinality</td>
<td>int</td>
<td>returns the total number of elements in the array, or 0 if the array is empty</td>
<td><code>cardinality(ARRAY[[1,2],[3,4]])</code></td>
<td></td>
</tr>
<tr>
<td>string_to_array</td>
<td>text[]</td>
<td>splits string into array elements using supplied delimiter and optional null string</td>
<td><code>string_to_array('xx~^~yy~^~zz', '~^~', 'yy')</code></td>
<td></td>
</tr>
<tr>
<td>unnest</td>
<td>setof</td>
<td>expand an array to a set of rows</td>
<td><code>unnest(ARRAY[1,2])</code></td>
<td>(2 rows)</td>
</tr>
</tbody>
</table>
### 9.19. Range Functions and Operators

See Section 8.17 for an overview of range types.

Table 9-49 shows the operators available for range types.

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>=</td>
<td>equal</td>
<td>int4range(1,5) = '1[4]'::int4range</td>
<td>t</td>
</tr>
<tr>
<td>&lt;&gt;</td>
<td>not equal</td>
<td>numrange(1.1,2.2) &lt;&gt; numrange(1.1,2.3)</td>
<td>t</td>
</tr>
<tr>
<td>&lt;</td>
<td>less than</td>
<td>int4range(1,10) &lt; int4range(2,3)</td>
<td>t</td>
</tr>
</tbody>
</table>

**Note:** There are two differences in the behavior of `string_to_array` from pre-9.1 versions of PostgreSQL. First, it will return an empty (zero-element) array rather than NULL when the input string is of zero length. Second, if the delimiter string is NULL, the function splits the input into individual characters, rather than returning NULL as before.

See also Section 9.20 about the aggregate function `array_agg` for use with arrays.
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>&gt;</td>
<td>greater than</td>
<td><code>int4range(1,10) &gt; int4range(1,5)</code></td>
<td>t</td>
</tr>
<tr>
<td>&lt;=</td>
<td>less than or equal</td>
<td><code>numrange(1.1,2.2) &lt;= numrange(1.1,2.2)</code></td>
<td>t</td>
</tr>
<tr>
<td>&gt;=</td>
<td>greater than or equal</td>
<td><code>numrange(1.1,2.2) &gt;= numrange(1.1,2.0)</code></td>
<td>t</td>
</tr>
<tr>
<td>@&gt;</td>
<td>contains range</td>
<td><code>int4range(2,4) @&gt; int4range(2,3)</code></td>
<td>t</td>
</tr>
<tr>
<td>@&gt;</td>
<td>contains element</td>
<td><code>'2011-01-01,2011-03-01'::tsrange @&gt; '2011-01-10'::timestamp</code></td>
<td>t</td>
</tr>
<tr>
<td>&lt;@</td>
<td>range is contained by</td>
<td><code>int4range(2,4) &lt;@ int4range(1,7)</code></td>
<td>t</td>
</tr>
<tr>
<td>&lt;@</td>
<td>element is contained by</td>
<td><code>42 &lt;@ int4range(1,7)</code></td>
<td>f</td>
</tr>
<tr>
<td>&amp;&amp;</td>
<td>overlap (have points in common)</td>
<td><code>int8range(3,7) &amp;&amp; int8range(4,12)</code></td>
<td>t</td>
</tr>
<tr>
<td>&lt;&lt;</td>
<td>strictly left of</td>
<td><code>int8range(1,10) &lt;&lt; int8range(100,110)</code></td>
<td>t</td>
</tr>
<tr>
<td>&gt;&gt;</td>
<td>strictly right of</td>
<td><code>int8range(50,60) &gt;&gt; int8range(20,30)</code></td>
<td>t</td>
</tr>
<tr>
<td>&amp;&lt;</td>
<td>does not extend to the right of</td>
<td><code>int8range(1,20) &amp;&lt; int8range(18,20)</code></td>
<td>t</td>
</tr>
<tr>
<td>&amp;&gt;</td>
<td>does not extend to the left of</td>
<td><code>int8range(7,20) &amp;&gt; int8range(5,10)</code></td>
<td>t</td>
</tr>
<tr>
<td>-</td>
<td>-</td>
<td>is adjacent</td>
<td>`numrange(1.1,2.2) -</td>
</tr>
<tr>
<td>+</td>
<td>union</td>
<td><code>numrange(5,15) + numrange(10,20)</code></td>
<td><code>{5,20}</code></td>
</tr>
<tr>
<td>*</td>
<td>intersection</td>
<td><code>int8range(5,15) * int8range(10,20)</code></td>
<td><code>{10,15}</code></td>
</tr>
<tr>
<td>-</td>
<td>difference</td>
<td><code>int8range(5,15) - int8range(10,20)</code></td>
<td><code>{5,10}</code></td>
</tr>
</tbody>
</table>

The simple comparison operators `<`, `>`, `<=`, and `>=` compare the lower bounds first, and only if those are equal, compare the upper bounds. These comparisons are not usually very useful for ranges, but are provided to allow B-tree indexes to be constructed on ranges.

The left-of/right-of/adjacent operators always return false when an empty range is involved; that is,
an empty range is not considered to be either before or after any other range.

The union and difference operators will fail if the resulting range would need to contain two disjoint sub-ranges, as such a range cannot be represented.

Table 9-50 shows the functions available for use with range types.

**Table 9-50. Range Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>lower(anyrange)</td>
<td>range’s element type</td>
<td>lower bound of range</td>
<td>lower(numrange(1.1,2.2))</td>
<td>1.1</td>
</tr>
<tr>
<td>upper(anyrange)</td>
<td>range’s element type</td>
<td>upper bound of range</td>
<td>upper(numrange(2.2,2.2))</td>
<td>2.2</td>
</tr>
<tr>
<td>isempty(anyrange)</td>
<td>boolean</td>
<td>is the range empty?</td>
<td>isempty(numrange(1.1,2.2))</td>
<td>false</td>
</tr>
<tr>
<td>lower_inc(anyrange)</td>
<td>boolean</td>
<td>is the lower bound inclusive?</td>
<td>lower_inc(numrange(1.1,2.2))</td>
<td>true</td>
</tr>
<tr>
<td>upper_inc(anyrange)</td>
<td>boolean</td>
<td>is the upper bound inclusive?</td>
<td>upper_inc(numrange(1.1,2.2))</td>
<td>true</td>
</tr>
<tr>
<td>lower_inf(anyrange)</td>
<td>boolean</td>
<td>is the lower bound infinite?</td>
<td>lower_inf(‘(,)’::daterange)</td>
<td>t</td>
</tr>
<tr>
<td>upper_inf(anyrange)</td>
<td>boolean</td>
<td>is the upper bound infinite?</td>
<td>upper_inf(‘(,)’::daterange)</td>
<td>t</td>
</tr>
<tr>
<td>range_merge(anyrange, anyrange)</td>
<td>anyrange</td>
<td>the smallest range which includes both of the given ranges</td>
<td>range_merge(’[1,2)’::int4range, ’(3,4)’::int4range)</td>
<td>[1,4)</td>
</tr>
</tbody>
</table>

The `lower` and `upper` functions return null if the range is empty or the requested bound is infinite. The `lower_inc`, `upper_inc`, `lower_inf`, and `upper_inf` functions all return false for an empty range.

### 9.20. Aggregate Functions

*Aggregate functions* compute a single result from a set of input values. The built-in normal aggregate functions are listed in Table 9-51 and Table 9-52. The built-in ordered-set aggregate functions are listed in Table 9-53 and Table 9-54. Grouping operations, which are closely related to aggregate functions, are listed in Table 9-55. The special syntax considerations for aggregate functions are explained in Section 4.2.7. Consult Section 2.7 for additional introductory information.

**Table 9-51. General-Purpose Aggregate Functions**


<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>array_agg(expression)</td>
<td>any non-array type</td>
<td>array of the argument type</td>
<td>No</td>
<td>input values, including nulls, concatenated into an array</td>
</tr>
<tr>
<td>array_agg(expression)</td>
<td>any array type</td>
<td>same as argument data type</td>
<td>No</td>
<td>input arrays concatenated into array of one higher dimension (inputs must all have same dimensionality, and cannot be empty or null)</td>
</tr>
<tr>
<td>avg(expression)</td>
<td>smallint, int, bigint, real, double precision, numeric, or interval</td>
<td>numeric for any integer-type argument, double precision for a floating-point argument, otherwise the same as the argument data type</td>
<td>Yes</td>
<td>the average (arithmetic mean) of all non-null input values</td>
</tr>
<tr>
<td>bit_and(expression)</td>
<td>smallint, int, bigint, or bit</td>
<td>same as argument data type</td>
<td>Yes</td>
<td>the bitwise AND of all non-null input values, or null if none</td>
</tr>
<tr>
<td>bit_or(expression)</td>
<td>smallint, int, bigint, or bit</td>
<td>same as argument data type</td>
<td>Yes</td>
<td>the bitwise OR of all non-null input values, or null if none</td>
</tr>
<tr>
<td>bool_and(expression)</td>
<td>bool</td>
<td>bool</td>
<td>Yes</td>
<td>true if all input values are true, otherwise false</td>
</tr>
<tr>
<td>bool_or(expression)</td>
<td>bool</td>
<td>bool</td>
<td>Yes</td>
<td>true if at least one input value is true, otherwise false</td>
</tr>
<tr>
<td>count(*)</td>
<td>bigint</td>
<td>bigint</td>
<td>Yes</td>
<td>number of input rows</td>
</tr>
<tr>
<td>count(expression)</td>
<td>any</td>
<td>bigint</td>
<td>Yes</td>
<td>number of input rows for which the value of expression is not null</td>
</tr>
<tr>
<td>Function</td>
<td>Argument Type(s)</td>
<td>Return Type</td>
<td>Partial Mode</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------</td>
<td>------------------</td>
<td>----------------</td>
<td>--------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>every(expression)</td>
<td>bool</td>
<td>bool</td>
<td>Yes</td>
<td>equivalent to bool_and</td>
</tr>
<tr>
<td>json_agg(expression)</td>
<td>any</td>
<td>json</td>
<td>No</td>
<td>aggregates values, including nulls, as a JSON array</td>
</tr>
<tr>
<td>jsonb_agg(expression)</td>
<td>any</td>
<td>jsonb</td>
<td>No</td>
<td>aggregates values, including nulls, as a JSON array</td>
</tr>
<tr>
<td>json_object_agg(name, value)</td>
<td>(any, any)</td>
<td>json</td>
<td>No</td>
<td>aggregates name/value pairs as a JSON object; values can be null, but not names</td>
</tr>
<tr>
<td>jsonb_object_agg(name, value)</td>
<td>(any, any)</td>
<td>jsonb</td>
<td>No</td>
<td>aggregates name/value pairs as a JSON object; values can be null, but not names</td>
</tr>
<tr>
<td>max(expression)</td>
<td>any numeric, string, date/time, network, or enum type, or arrays of these types</td>
<td>same as argument type</td>
<td>Yes</td>
<td>maximum value of expression across all non-null input values</td>
</tr>
<tr>
<td>min(expression)</td>
<td>any numeric, string, date/time, network, or enum type, or arrays of these types</td>
<td>same as argument type</td>
<td>Yes</td>
<td>minimum value of expression across all non-null input values</td>
</tr>
<tr>
<td>string_agg(expression, delimiter)</td>
<td>(text, text) or (text, bytea)</td>
<td>same as argument types</td>
<td>No</td>
<td>non-null input values concatenated into a string, separated by delimiter</td>
</tr>
<tr>
<td>sum(expression)</td>
<td>smallint, int, bigint, real, double precision, numeric, interval, or money</td>
<td>bigint for smallint or int arguments, numeric for bigint arguments, otherwise the same as the argument data type</td>
<td>Yes</td>
<td>sum of expression across all non-null input values</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>xmlagg(expression)</td>
<td>xml</td>
<td>xml</td>
<td>No</td>
<td>concatenation of non-null XML values (see also Section 9.14.1.7)</td>
</tr>
</tbody>
</table>

It should be noted that except for `count`, these functions return a null value when no rows are selected. In particular, `sum` of no rows returns null, not zero as one might expect, and `array_agg` returns null rather than an empty array when there are no input rows. The `coalesce` function can be used to substitute zero or an empty array for null when necessary.

Aggregate functions which support *Partial Mode* are eligible to participate in various optimizations, such as parallel aggregation.

**Note:** Boolean aggregates `bool_and` and `bool_or` correspond to standard SQL aggregates `every` and `any` or `some`. As for `any` and `some`, it seems that there is an ambiguity built into the standard syntax:

```sql
SELECT b1 = ANY((SELECT b2 FROM t2 ...)) FROM t1 ...;
```

Here `ANY` can be considered either as introducing a subquery, or as being an aggregate function, if the subquery returns one row with a Boolean value. Thus the standard name cannot be given to these aggregates.

**Note:** Users accustomed to working with other SQL database management systems might be disappointed by the performance of the `count` aggregate when it is applied to the entire table. A query like:

```sql
SELECT count(*) FROM sometable;
```

will require effort proportional to the size of the table: PostgreSQL will need to scan either the entire table or the entirety of an index which includes all rows in the table.

The aggregate functions `array_agg`, `json_agg`, `jsonb_agg`, `json_object_agg`, `jsonb_object_agg`, `string_agg`, and `xmlagg`, as well as similar user-defined aggregate functions, produce meaningfully different result values depending on the order of the input values. This ordering is unspecified by default, but can be controlled by writing an `ORDER BY` clause within the aggregate call, as shown in Section 4.2.7. Alternatively, supplying the input values from a sorted subquery will usually work. For example:

```sql
SELECT xmlagg(x) FROM (SELECT x FROM test ORDER BY y DESC) AS tab;
```

Beware that this approach can fail if the outer query level contains additional processing, such as a join, because that might cause the subquery’s output to be reordered before the aggregate is computed.

Table 9-52 shows aggregate functions typically used in statistical analysis. (These are separated out merely to avoid cluttering the listing of more-commonly-used aggregates.) Where the description mentions $N$, it means the number of input rows for which all the input expressions are non-null. In all cases, null is returned if the computation is meaningless, for example when $N$ is zero.

**Table 9-52. Aggregate Functions for Statistics**
<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>corr(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>correlation coefficient</td>
</tr>
<tr>
<td>covar_pop(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>population covariance</td>
</tr>
<tr>
<td>covar_samp(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>sample covariance</td>
</tr>
<tr>
<td>regr_avgx(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>average of the independent variable ( \frac{\sum Y}{N} )</td>
</tr>
<tr>
<td>regr_avgy(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>average of the dependent variable ( \frac{\sum Y}{N} )</td>
</tr>
<tr>
<td>regr_count(Y, X)</td>
<td>double precision</td>
<td>bigint</td>
<td>Yes</td>
<td>number of input rows in which both expressions are nonnull</td>
</tr>
<tr>
<td>regr_intercept(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>y-intercept of the least-squares-fit linear equation determined by the ((X, Y)) pairs</td>
</tr>
<tr>
<td>regr_r2(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>square of the correlation coefficient</td>
</tr>
<tr>
<td>regr_slope(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>slope of the least-squares-fit linear equation determined by the ((X, Y)) pairs</td>
</tr>
<tr>
<td>regr_sxx(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>( \sum (X^2) - \frac{(\sum X)^2}{N} ) \text{ \textquotedblleft sum of squares\textquotedblright} of the independent variable)</td>
</tr>
<tr>
<td>regr_sxy(Y, X)</td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>( \sum (X \cdot Y) - \frac{\sum X \cdot \sum Y}{N} ) \text{ \textquotedblleft sum of products\textquotedblright of independent times dependent variable)</td>
</tr>
</tbody>
</table>
Table 9-53 shows some aggregate functions that use the \textit{ordered-set aggregate} syntax. These functions are sometimes referred to as “inverse distribution” functions.

### Table 9-53. Ordered-Set Aggregate Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>regr_syy(Y, X)</code></td>
<td>double precision</td>
<td>double precision</td>
<td>Yes</td>
<td>$\text{sum}(Y^2) - \frac{\text{sum}(Y)^2}{N}$ (&quot;sum of squares&quot; of the dependent variable)</td>
</tr>
<tr>
<td><code>stddev(expression)</code></td>
<td>smallint, int, bigint, real, double precision, or numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>historical alias for <code>stddev_samp</code></td>
</tr>
<tr>
<td><code>stddev_pop(expression)</code></td>
<td>smallint, int, bigint, real, double precision, or numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>population standard deviation of the input values</td>
</tr>
<tr>
<td><code>stddev_samp(expression)</code></td>
<td>smallint, int, bigint, real, double precision, or numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>sample standard deviation of the input values</td>
</tr>
<tr>
<td><code>variance(expression)</code></td>
<td>smallint, int, bigint, real, double precision, or numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>historical alias for <code>var_samp</code></td>
</tr>
<tr>
<td><code>var_pop(expression)</code></td>
<td>smallint, int, bigint, real, double precision, or numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>population variance of the input values (square of the population standard deviation)</td>
</tr>
<tr>
<td><code>var_samp(expression)</code></td>
<td>smallint, int, bigint, real, double precision, or numeric</td>
<td>double precision</td>
<td>Yes</td>
<td>sample variance of the input values (square of the sample standard deviation)</td>
</tr>
<tr>
<td>Function</td>
<td>Direct Argument Type(s)</td>
<td>Aggregated Argument Type(s)</td>
<td>Return Type</td>
<td>Partial Mode</td>
</tr>
<tr>
<td>----------</td>
<td>-------------------------</td>
<td>-----------------------------</td>
<td>-------------</td>
<td>--------------</td>
</tr>
<tr>
<td>mode() WITHIN GROUP (ORDER BY sort_expression)</td>
<td>any sortable type</td>
<td>same as sort expression</td>
<td>No</td>
<td>returns the most frequent input value (arbitrarily choosing the first one if there are multiple equally-frequent results)</td>
</tr>
<tr>
<td>percentile_cont(fraction) WITHIN GROUP (ORDER BY sort_expression)</td>
<td>double precision</td>
<td>double precision or interval</td>
<td>same as sort expression</td>
<td>No</td>
</tr>
<tr>
<td>percentile_cont(fractions) WITHIN GROUP (ORDER BY sort_expression)</td>
<td>double precision[]</td>
<td>array of sort expression’s type</td>
<td>No</td>
<td>multiple continuous percentile: returns an array of results matching the shape of the fractions parameter, with each non-null element replaced by the value corresponding to that percentile</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

#### Table 9-53. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Direct Argument Type(s)</th>
<th>Aggregated Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>percentile_disc(fraction) WITHIN GROUP (ORDER BY sort_expression)</td>
<td>double precision</td>
<td>any sortable type</td>
<td>same as sort expression</td>
<td>No</td>
<td>discrete percentile: returns the first input value whose position in the ordering equals or exceeds the specified fraction</td>
</tr>
<tr>
<td>percentile_disc(fractions) WITHIN GROUP (ORDER BY sort_expression)</td>
<td>double precision[]</td>
<td>any sortable type</td>
<td>array of sort expression’s type</td>
<td>No</td>
<td>multiple discrete percentile: returns an array of results matching the shape of the fractions parameter, with each non-null element replaced by the input value corresponding to that percentile</td>
</tr>
</tbody>
</table>

All the aggregates listed in Table 9-53 ignore null values in their sorted input. For those that take a fraction parameter, the fraction value must be between 0 and 1; an error is thrown if not. However, a null fraction value simply produces a null result.

Each of the aggregates listed in Table 9-54 is associated with a window function of the same name defined in Section 9.21. In each case, the aggregate result is the value that the associated window function would have returned for the “hypothetical” row constructed from args, if such a row had been added to the sorted group of rows computed from the sorted_args.

#### Table 9-54. Hypothetical-Set Aggregate Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Direct Argument Type(s)</th>
<th>Aggregated Argument Type(s)</th>
<th>Return Type</th>
<th>Partial Mode</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rank(args) WITHIN GROUP (ORDER BY sorted_args)</td>
<td>VARIADIC &quot;any&quot;</td>
<td>VARIADIC &quot;any&quot;</td>
<td>bigint</td>
<td>No</td>
<td>rank of the hypothetical row, with gaps for duplicate rows</td>
</tr>
</tbody>
</table>
For each of these hypothetical-set aggregates, the list of direct arguments given in \texttt{args} must match the number and types of the aggregated arguments given in \texttt{sorted_args}. Unlike most built-in aggregates, these aggregates are not strict, that is they do not drop input rows containing nulls. Null values sort according to the rule specified in the \texttt{ORDER BY} clause.

**Table 9-55. Grouping Operations**

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{GROUPING(args...)}</td>
<td>integer</td>
<td>Integer bit mask indicating which arguments are not being included in the current grouping set</td>
</tr>
</tbody>
</table>

Grouping operations are used in conjunction with grouping sets (see Section 7.2.4) to distinguish result rows. The arguments to the \texttt{GROUPING} operation are not actually evaluated, but they must match exactly expressions given in the \texttt{GROUP BY} clause of the associated query level. Bits are assigned with the rightmost argument being the least-significant bit; each bit is 0 if the corresponding expression is included in the grouping criteria of the grouping set generating the result row, and 1 if it is not. For example:

```
=> SELECT * FROM items_sold;
make | model | sales
-------+-------+-------
Foo   | GT    | 10
Foo   | Tour  | 20
Bar   | City  | 15
Bar   | Sport | 5
```
9.21. Window Functions

Window functions provide the ability to perform calculations across sets of rows that are related to the current query row. See Section 3.5 for an introduction to this feature, and Section 4.2.8 for syntax details.

The built-in window functions are listed in Table 9-56. Note that these functions must be invoked using window function syntax; that is an OVER clause is required.

In addition to these functions, any built-in or user-defined normal aggregate function (but not ordered-set or hypothetical-set aggregates) can be used as a window function; see Section 9.20 for a list of the built-in aggregates. Aggregate functions act as window functions only when an OVER clause follows the call; otherwise they act as regular aggregates.

Table 9-56. General-Purpose Window Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>row_number()</td>
<td>bigint</td>
<td>number of the current row within its partition, counting from 1</td>
</tr>
<tr>
<td>rank()</td>
<td>bigint</td>
<td>rank of the current row with gaps; same as row_number of its first peer</td>
</tr>
<tr>
<td>dense_rank()</td>
<td>bigint</td>
<td>rank of the current row without gaps; this function counts peer groups</td>
</tr>
<tr>
<td>percent_rank()</td>
<td>double precision</td>
<td>relative rank of the current row: (rank - 1) / (total rows - 1)</td>
</tr>
<tr>
<td>cume_dist()</td>
<td>double precision</td>
<td>relative rank of the current row: (number of rows preceding or peer with current row) / (total rows)</td>
</tr>
<tr>
<td>ntile(num_buckets integer)</td>
<td>integer</td>
<td>integer ranging from 1 to the argument value, dividing the partition as equally as possible</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>lag(value anyelement [, offset integer [, default anyelement ]])</code></td>
<td>same type as value</td>
<td>returns value evaluated at the row that is offset rows before the current row within the partition; if there is no such row, instead return default (which must be of the same type as value). Both offset and default are evaluated with respect to the current row. If omitted, offset defaults to 1 and default to null.</td>
</tr>
<tr>
<td><code>lead(value anyelement [, offset integer [, default anyelement ]])</code></td>
<td>same type as value</td>
<td>returns value evaluated at the row that is offset rows after the current row within the partition; if there is no such row, instead return default (which must be of the same type as value). Both offset and default are evaluated with respect to the current row. If omitted, offset defaults to 1 and default to null.</td>
</tr>
<tr>
<td><code>first_value(value any)</code></td>
<td>same type as value</td>
<td>returns value evaluated at the row that is the first row of the window frame</td>
</tr>
<tr>
<td><code>last_value(value any)</code></td>
<td>same type as value</td>
<td>returns value evaluated at the row that is the last row of the window frame</td>
</tr>
<tr>
<td><code>nth_value(value any, nth integer)</code></td>
<td>same type as value</td>
<td>returns value evaluated at the row that is the nth row of the window frame (counting from 1); null if no such row</td>
</tr>
</tbody>
</table>

All of the functions listed in Table 9-56 depend on the sort ordering specified by the `ORDER BY` clause of the associated window definition. Rows that are not distinct in the `ORDER BY` ordering are said to be peers; the four ranking functions are defined so that they give the same answer for any two peer rows.

Note that `first_value`, `last_value`, and `nth_value` consider only the rows within the “window frame”, which by default contains the rows from the start of the partition through the last peer of the current row. This is likely to give unhelpful results for `last_value` and sometimes also `nth_value`. You can redefine the frame by adding a suitable frame specification (`RANGE` or `ROWS`) to the `OVER` clause. See Section 4.2.8 for more information about frame specifications.

When an aggregate function is used as a window function, it aggregates over the rows within the current row’s window frame. An aggregate used with `ORDER BY` and the default window frame definition produces a “running sum” type of behavior, which may or may not be what’s wanted. To obtain aggregation over the whole partition, omit `ORDER BY` or use `ROWS BETWEEN UNBOUNDED PRECEDING AND UNBOUNDED FOLLOWING`. Other frame specifications can be used to obtain other effects.

**Note:** The SQL standard defines a `RESPECT NULLS` or `IGNORE NULLS` option for `lead`, `lag`, 

300
first_value, last_value, and nth_value. This is not implemented in PostgreSQL; the behavior is always the same as the standard’s default, namely RESPECT NULLS. Likewise, the standard’s FROM FIRST or FROM LAST option for nth_value is not implemented; only the default FROM FIRST behavior is supported. (You can achieve the result of FROM LAST by reversing the ORDER BY ordering.)

9.22. Subquery Expressions

This section describes the SQL-compliant subquery expressions available in PostgreSQL. All of the expression forms documented in this section return Boolean (true/false) results.

9.22.1. EXISTS

EXISTS (subquery)

The argument of EXISTS is an arbitrary SELECT statement, or subquery. The subquery is evaluated to determine whether it returns any rows. If it returns at least one row, the result of EXISTS is “true”; if the subquery returns no rows, the result of EXISTS is “false”.

The subquery can refer to variables from the surrounding query, which will act as constants during any one evaluation of the subquery.

The subquery will generally only be executed long enough to determine whether at least one row is returned, not all the way to completion. It is unwise to write a subquery that has side effects (such as calling sequence functions); whether the side effects occur might be unpredictable.

Since the result depends only on whether any rows are returned, and not on the contents of those rows, the output list of the subquery is normally unimportant. A common coding convention is to write all EXISTS tests in the form EXISTS(SELECT 1 WHERE ...). There are exceptions to this rule however, such as subqueries that use INTERSECT.

This simple example is like an inner join on col2, but it produces at most one output row for each tab1 row, even if there are several matching tab2 rows:

```sql
SELECT col1
FROM tab1
WHERE EXISTS (SELECT 1 FROM tab2 WHERE col2 = tab1.col2);
```

9.22.2. IN

expression IN (subquery)

The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result. The result of IN is “true” if any equal subquery row is found. The result is “false” if no equal row is found (including the case where the subquery returns no rows).
Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand row yields null, the result of the `IN` construct will be null, not false. This is in accordance with SQL’s normal rules for Boolean combinations of null values.

As with `EXISTS`, it’s unwise to assume that the subquery will be evaluated completely.

```
row_constructor IN (subquery)
```

The left-hand side of this form of `IN` is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result. The result of `IN` is “true” if any equal subquery row is found. The result is “false” if no equal row is found (including the case where the subquery returns no rows).

As usual, null values in the rows are combined per the normal rules of SQL Boolean expressions. Two rows are considered equal if all their corresponding members are non-null and equal; the rows are unequal if any corresponding members are non-null and unequal; otherwise the result of that row comparison is unknown (null). If all the per-row results are either unequal or null, with at least one null, then the result of `IN` is null.

### 9.22.3. NOT IN

```
expression NOT IN (subquery)
```

The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result. The result of `NOT IN` is “true” if only unequal subquery rows are found (including the case where the subquery returns no rows). The result is “false” if any equal row is found.

Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand row yields null, the result of the `NOT IN` construct will be null, not true. This is in accordance with SQL’s normal rules for Boolean combinations of null values.

As with `EXISTS`, it’s unwise to assume that the subquery will be evaluated completely.

```
row_constructor NOT IN (subquery)
```

The left-hand side of this form of `NOT IN` is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result. The result of `NOT IN` is “true” if only unequal subquery rows are found (including the case where the subquery returns no rows). The result is “false” if any equal row is found.

As usual, null values in the rows are combined per the normal rules of SQL Boolean expressions. Two rows are considered equal if all their corresponding members are non-null and equal; the rows are unequal if any corresponding members are non-null and unequal; otherwise the result of that row comparison is unknown (null). If all the per-row results are either unequal or null, with at least one null, then the result of `NOT IN` is null.

### 9.22.4. ANY/SOME

```
eexpression operator ANY (subquery)
eexpression operator SOME (subquery)
```
Chapter 9. Functions and Operators

The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result using the given operator, which must yield a Boolean result. The result of \textit{ANY} is “true” if any true result is obtained. The result is “false” if no true result is found (including the case where the subquery returns no rows).

\texttt{SOME} is a synonym for \textit{ANY}. \texttt{IN} is equivalent to \texttt{= ANY}.

Note that if there are no successes and at least one right-hand row yields null for the operator’s result, the result of the \textit{ANY} construct will be null, not false. This is in accordance with SQL’s normal rules for Boolean combinations of null values.

As with \texttt{EXISTS}, it’s unwise to assume that the subquery will be evaluated completely.

\begin{verbatim}
row_constructor operator ANY (subquery)
row_constructor operator SOME (subquery)
\end{verbatim}

The left-hand side of this form of \textit{ANY} is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result, using the given operator. The result of \textit{ANY} is “true” if the comparison returns true for any subquery row. The result is “false” if the comparison returns false for every subquery row (including the case where the subquery returns no rows). The result is NULL if no comparison with a subquery row returns true, and at least one comparison returns NULL.

See Section 9.23.5 for details about the meaning of a row constructor comparison.

\subsection*{9.22.5. \texttt{ALL}}

\begin{verbatim}
expression operator ALL (subquery)
\end{verbatim}

The right-hand side is a parenthesized subquery, which must return exactly one column. The left-hand expression is evaluated and compared to each row of the subquery result using the given operator, which must yield a Boolean result. The result of \textit{ALL} is “true” if all rows yield true (including the case where the subquery returns no rows). The result is “false” if any false result is found. The result is NULL if no comparison with a subquery row returns false, and at least one comparison returns NULL.

\texttt{NOT IN} is equivalent to \texttt{<> ALL}.

As with \texttt{EXISTS}, it’s unwise to assume that the subquery will be evaluated completely.

\begin{verbatim}
row_constructor operator ALL (subquery)
\end{verbatim}

The left-hand side of this form of \textit{ALL} is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. The left-hand expressions are evaluated and compared row-wise to each row of the subquery result, using the given operator. The result of \textit{ALL} is “true” if the comparison returns true for all subquery rows (including the case where the subquery returns no rows). The result is “false” if the comparison returns false for any subquery row. The result is NULL if no comparison with a subquery row returns false, and at least one comparison returns NULL.

See Section 9.23.5 for details about the meaning of a row constructor comparison.
9.22.6. Single-row Comparison

`row_constructor operator (subquery)`

The left-hand side is a row constructor, as described in Section 4.2.13. The right-hand side is a parenthesized subquery, which must return exactly as many columns as there are expressions in the left-hand row. Furthermore, the subquery cannot return more than one row. (If it returns zero rows, the result is taken to be null.) The left-hand side is evaluated and compared row-wise to the single subquery result row.

See Section 9.23.5 for details about the meaning of a row constructor comparison.

9.23. Row and Array Comparisons

This section describes several specialized constructs for making multiple comparisons between groups of values. These forms are syntactically related to the subquery forms of the previous section, but do not involve subqueries. The forms involving array subexpressions are PostgreSQL extensions; the rest are SQL-compliant. All of the expression forms documented in this section return Boolean (true/false) results.

9.23.1. IN

`expression IN (value [, ...])`

The right-hand side is a parenthesized list of scalar expressions. The result is “true” if the left-hand expression’s result is equal to any of the right-hand expressions. This is a shorthand notation for

```
expression = value1
OR
expression = value2
OR
...
```

Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand expression yields null, the result of the IN construct will be null, not false. This is in accordance with SQL’s normal rules for Boolean combinations of null values.

9.23.2. NOT IN

`expression NOT IN (value [, ...])`

The right-hand side is a parenthesized list of scalar expressions. The result is “true” if the left-hand expression’s result is unequal to all of the right-hand expressions. This is a shorthand notation for

```
expression <> value1
AND
expression <> value2
AND
...
```
Chapter 9. Functions and Operators

Note that if the left-hand expression yields null, or if there are no equal right-hand values and at least one right-hand expression yields null, the result of the \texttt{NOT IN} construct will be null, not true as one might naively expect. This is in accordance with SQL’s normal rules for Boolean combinations of null values.

\textbf{Tip:} \( x \ NOT \ IN \ y \) is equivalent to \( \texttt{NOT (x \ IN \ y)} \) in all cases. However, null values are much more likely to trip up the novice when working with \( \texttt{NOT IN} \) than when working with \( \texttt{IN} \). It is best to express your condition positively if possible.

9.23.3. \textsc{any/some} (array)

\texttt{expression operator ANY (array expression)}
\texttt{expression operator SOME (array expression)}

The right-hand side is a parenthesized expression, which must yield an array value. The left-hand expression is evaluated and compared to each element of the array using the given \texttt{operator}, which must yield a Boolean result. The result of \texttt{ANY} is “true” if any true result is obtained. The result is “false” if no true result is found (including the case where the array has zero elements).

If the array expression yields a null array, the result of \texttt{ANY} will be null. If the left-hand expression yields null, the result of \texttt{ANY} is ordinarily null (though a non-strict comparison operator could possibly yield a different result). Also, if the right-hand array contains any null elements and no true comparison result is obtained, the result of \texttt{ANY} will be null, not false (again, assuming a strict comparison operator). This is in accordance with SQL’s normal rules for Boolean combinations of null values.

\texttt{SOME} is a synonym for \texttt{ANY}.

9.23.4. \textsc{all} (array)

\texttt{expression operator ALL (array expression)}

The right-hand side is a parenthesized expression, which must yield an array value. The left-hand expression is evaluated and compared to each element of the array using the given \texttt{operator}, which must yield a Boolean result. The result of \texttt{ALL} is “true” if all comparisons yield true (including the case where the array has zero elements). The result is “false” if any false result is found.

If the array expression yields a null array, the result of \texttt{ALL} will be null. If the left-hand expression yields null, the result of \texttt{ALL} is ordinarily null (though a non-strict comparison operator could possibly yield a different result). Also, if the right-hand array contains any null elements and no false comparison result is obtained, the result of \texttt{ALL} will be null, not true (again, assuming a strict comparison operator). This is in accordance with SQL’s normal rules for Boolean combinations of null values.

9.23.5. Row Constructor Comparison

\texttt{row_constructor operator row_constructor}

Each side is a row constructor, as described in Section 4.2.13. The two row values must have the same number of fields. Each side is evaluated and they are compared row-wise. Row constructor
comparisons are allowed when the operator is =, <>, <, <=, > or >=. Every row element must be of a type which has a default B-tree operator class or the attempted comparison may generate an error.

**Note:** Errors related to the number or types of elements might not occur if the comparison is resolved using earlier columns.

The = and < > cases work slightly differently from the others. Two rows are considered equal if all their corresponding members are non-null and equal; the rows are unequal if any corresponding members are non-null and unequal; otherwise the result of the row comparison is unknown (null).

For the <, <=, > and > = cases, the row elements are compared left-to-right, stopping as soon as an unequal or null pair of elements is found. If either of this pair of elements is null, the result of the row comparison is unknown (null); otherwise comparison of this pair of elements determines the result. For example, ROW (1, 2, NULL) < ROW (1, 3, 0) yields true, not null, because the third pair of elements are not considered.

**Note:** Prior to PostgreSQL 8.2, the <, <=, > and > = cases were not handled per SQL specification. A comparison like ROW (a, b) < ROW (c, d) was implemented as a < c AND b < d whereas the correct behavior is equivalent to a < c OR (a = c AND b < d).

row_constructor IS DISTINCT FROM row_constructor

This construct is similar to a < > row comparison, but it does not yield null for null inputs. Instead, any null value is considered unequal to (distinct from) any non-null value, and any two nulls are considered equal (not distinct). Thus the result will either be true or false, never null.

row_constructor IS NOT DISTINCT FROM row_constructor

This construct is similar to a = row comparison, but it does not yield null for null inputs. Instead, any null value is considered unequal to (distinct from) any non-null value, and any two nulls are considered equal (not distinct). Thus the result will always be either true or false, never null.

### 9.23.6. Composite Type Comparison

record operator record

The SQL specification requires row-wise comparison to return NULL if the result depends on comparing two NULL values or a NULL and a non-NULL. PostgreSQL does this only when comparing the results of two row constructors (as in Section 9.23.5) or comparing a row constructor to the output of a subquery (as in Section 9.22). In other contexts where two composite-type values are compared, two NULL field values are considered equal, and a NULL is considered larger than a non-NULL. This is necessary in order to have consistent sorting and indexing behavior for composite types.

Each side is evaluated and they are compared row-wise. Composite type comparisons are allowed when the operator is =, <>, <, <=, > or >=. or has semantics similar to one of these. (To be specific, an operator can be a row comparison operator if it is a member of a B-tree operator class, or is the negator of the = member of a B-tree operator class.) The default behavior of the above operators is the same as for IS [ NOT ] DISTINCT FROM for row constructors (see Section 9.23.5).
To support matching of rows which include elements without a default B-tree operator class, the following operators are defined for composite type comparison: \( \ast\leq, \ast<, \ast<, \ast\leq, \ast>, \) and \( \ast\geq\).

These operators compare the internal binary representation of the two rows. Two rows might have a different binary representation even though comparisons of the two rows with the equality operator is true. The ordering of rows under these comparison operators is deterministic but not otherwise meaningful. These operators are used internally for materialized views and might be useful for other specialized purposes such as replication but are not intended to be generally useful for writing queries.

### 9.24. Set Returning Functions

This section describes functions that possibly return more than one row. The most widely used functions in this class are series generating functions, as detailed in Table 9-57 and Table 9-58. Other, more specialized set-returning functions are described elsewhere in this manual. See Section 7.2.1.4 for ways to combine multiple set-returning functions.

#### Table 9-57. Series Generating Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Argument Type</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>generate_series(start, stop)</code></td>
<td>int, bigint or numeric</td>
<td>setof int, setof bigint, or setof numeric (same as argument type)</td>
<td>Generate a series of values, from start to stop with a step size of one</td>
</tr>
<tr>
<td><code>generate_series(start, stop, step)</code></td>
<td>int, bigint or numeric</td>
<td>setof int, setof bigint or setof numeric (same as argument type)</td>
<td>Generate a series of values, from start to stop with a step size of step</td>
</tr>
<tr>
<td><code>generate_series(start, stop, step interval)</code></td>
<td>timestamp or timestamp with time zone</td>
<td>setof timestamp or setof timestamp with time zone (same as argument type)</td>
<td>Generate a series of values, from start to stop with a step size of step</td>
</tr>
</tbody>
</table>

When `step` is positive, zero rows are returned if `start` is greater than `stop`. Conversely, when `step` is negative, zero rows are returned if `start` is less than `stop`. Zero rows are also returned for `NULL` inputs. It is an error for `step` to be zero. Some examples follow:

```sql
SELECT * FROM generate_series(2,4);
generate_series
-----------------
  2
  3
  4
(3 rows)
```

```sql
SELECT * FROM generate_series(5,1,-2);
generate_series
-----------------
  5
  3
  1
(3 rows)
```
SELECT * FROM generate_series(4,3);
generate_series
-----------------
(0 rows)

SELECT generate_series(1.1, 4, 1.3);
generate_series
-----------------
1.1
2.4
3.7
(3 rows)

-- this example relies on the date-plus-integer operator
SELECT current_date + s.a AS dates FROM generate_series(0,14,7) AS s(a);
dates
------------
2004-02-05
2004-02-12
2004-02-19
(3 rows)

SELECT * FROM generate_series('2008-03-01 00:00'::timestamp, '2008-03-04 12:00', '10 hours');
generate_series
---------------------
2008-03-01 00:00:00
2008-03-01 10:00:00
2008-03-01 20:00:00
2008-03-02 06:00:00
2008-03-02 16:00:00
2008-03-03 02:00:00
2008-03-03 12:00:00
2008-03-03 22:00:00
2008-03-04 08:00:00
(9 rows)

Table 9-58. Subscript Generating Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>generate_subscripts(array anyarray, dim int)</td>
<td>setof int</td>
<td>Generate a series comprising the given array’s subscripts.</td>
</tr>
<tr>
<td>generate_subscripts(array anyarray, dim int, reverse boolean)</td>
<td>setof int</td>
<td>Generate a series comprising the given array’s subscripts. When reverse is true, the series is returned in reverse order.</td>
</tr>
</tbody>
</table>

`generate_subscripts` is a convenience function that generates the set of valid subscripts for the specified dimension of the given array. Zero rows are returned for arrays that do not have the requested dimension, or for NULL arrays (but valid subscripts are returned for NULL array elements). Some examples follow:
-- basic usage
SELECT generate_subscripts('{NULL,1,NULL,2}'::int[], 1) AS s;

---
1
2
3
4
(4 rows)

-- presenting an array, the subscript and the subscripted value requires a subquery
SELECT * FROM arrays;
a
--------------------
{-1,-2}
{100,200,300}
(2 rows)

SELECT a AS array, s AS subscript, a[s] AS value
FROM (SELECT generate_subscripts(a, 1) AS s, a FROM arrays) foo;

<table>
<thead>
<tr>
<th>array</th>
<th>subscript</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>{-1,-2}</td>
<td>1</td>
<td>-1</td>
</tr>
<tr>
<td>{-1,-2}</td>
<td>2</td>
<td>-2</td>
</tr>
<tr>
<td>{100,200,300}</td>
<td>1</td>
<td>100</td>
</tr>
<tr>
<td>{100,200,300}</td>
<td>2</td>
<td>200</td>
</tr>
<tr>
<td>{100,200,300}</td>
<td>3</td>
<td>300</td>
</tr>
</tbody>
</table>
(5 rows)

-- unnest a 2D array
CREATE OR REPLACE FUNCTION unnest2(anyarray)
RETURNS SETOF anyelement AS $$
select $1[i][j]
from generate_subscripts($1,1) g1(i),
generate_subscripts($1,2) g2(j);
$$ LANGUAGE sql IMMUTABLE;
CREATE FUNCTION

SELECT * FROM unnest2(ARRAY[[1,2],[3,4]]);
unnest2

---
1
2
3
4
(4 rows)

When a function in the FROM clause is suffixed by WITH ORDINALITY, a bigint column is appended to the output which starts from 1 and increments by 1 for each row of the function’s output. This is most useful in the case of set returning functions such as unnest().

-- set returning function WITH ORDINALITY
SELECT * FROM pg_ls_dir('.') WITH ORDINALITY AS t(ls,n);
tls  | n
-----------------
pg_serial | 1

309
9.25. System Information Functions

Table 9-59 shows several functions that extract session and system information.

In addition to the functions listed in this section, there are a number of functions related to the statistics system that also provide system information. See Section 28.2.2 for more information.

Table 9-59. Session Information Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>current_catalog</td>
<td>name</td>
<td>name of current database (called “catalog” in the SQL standard)</td>
</tr>
<tr>
<td>current_database()</td>
<td>name</td>
<td>name of current database</td>
</tr>
<tr>
<td>current_query()</td>
<td>text</td>
<td>text of the currently executing query, as submitted by the client (might contain more than one statement)</td>
</tr>
<tr>
<td>current_role</td>
<td>name</td>
<td>equivalent to current_user</td>
</tr>
<tr>
<td>current_schema[][]</td>
<td>name</td>
<td>name of current schema</td>
</tr>
<tr>
<td>current_schemas(boolean)</td>
<td>name[]</td>
<td>names of schemas in search path, optionally including implicit schemas</td>
</tr>
<tr>
<td>current_user</td>
<td>name</td>
<td>user name of current execution context</td>
</tr>
<tr>
<td>inet_client_addr()</td>
<td>inet</td>
<td>address of the remote connection</td>
</tr>
<tr>
<td>inet_client_port()</td>
<td>int</td>
<td>port of the remote connection</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>inet_server_addr()</td>
<td>inet</td>
<td>address of the local connection</td>
</tr>
<tr>
<td>inet_server_port()</td>
<td>int</td>
<td>port of the local connection</td>
</tr>
<tr>
<td>pg_backend_pid()</td>
<td>int</td>
<td>Process ID of the server process attached to the current session</td>
</tr>
<tr>
<td>pg_blocking_pids(int)</td>
<td>int[]</td>
<td>Process ID(s) that are blocking specified server process ID</td>
</tr>
<tr>
<td>pg_conf_load_time()</td>
<td>timestamp</td>
<td>configuration load time</td>
</tr>
<tr>
<td>pg_my_temp_schema()</td>
<td>oid</td>
<td>OID of session’s temporary schema, or 0 if none</td>
</tr>
<tr>
<td>pg_is_other_temp_schema(oid)</td>
<td>boolean</td>
<td>is schema another session’s temporary schema?</td>
</tr>
<tr>
<td>pg_listening_channels()</td>
<td>setof text</td>
<td>channel names that the session is currently listening on</td>
</tr>
<tr>
<td>pg_notification_queue_usage()</td>
<td>double</td>
<td>fraction of the asynchronous notification queue currently occupied (0-1)</td>
</tr>
<tr>
<td>pg_postmaster_start_time()</td>
<td>timestamp</td>
<td>server start time</td>
</tr>
<tr>
<td>pg_trigger_depth()</td>
<td>int</td>
<td>current nesting level of PostgreSQL triggers (0 if not called, directly or indirectly, from inside a trigger)</td>
</tr>
<tr>
<td>session_user</td>
<td>name</td>
<td>session user name</td>
</tr>
<tr>
<td>user</td>
<td>name</td>
<td>equivalent to current_user</td>
</tr>
<tr>
<td>version()</td>
<td>text</td>
<td>PostgreSQL version information. See also server_version_num for a machine-readable version.</td>
</tr>
</tbody>
</table>

**Note:** current_catalog, current_role, current_schema, current_user, session_user, and user have special syntactic status in SQL: they must be called without trailing parentheses. (In PostgreSQL, parentheses can optionally be used with current_schema, but not with the others.)

The session_user is normally the user who initiated the current database connection; but superusers can change this setting with SET SESSION AUTHORIZATION. The current_user is the user identifier that is applicable for permission checking. Normally it is equal to the session user, but it can be changed with SET ROLE. It also changes during the execution of functions with the attribute SECURITY DEFINER. In Unix parlance, the session user is the “real user” and the current user is the “effective user”. current_role and user are synonyms for current_user. (The SQL standard draws a distinction between current_role and current_user, but PostgreSQL does not, since it unifies users and roles into a single kind of entity.)

current_schema returns the name of the schema that is first in the search path (or a null value if the search path is empty). This is the schema that will be used for any tables or other named objects that are created without specifying a target schema. current_schemas(boolean) returns an array of
the names of all schemas presently in the search path. The Boolean option determines whether or not implicitly included system schemas such as \texttt{pg\_catalog} are included in the returned search path.

\textbf{Note:} The search path can be altered at run time. The command is:

\begin{verbatim}
SET search_path TO schema [, schema, ...]
\end{verbatim}

\texttt{inet\_client\_addr} returns the IP address of the current client, and \texttt{inet\_client\_port} returns the port number. \texttt{inet\_server\_addr} returns the IP address on which the server accepted the current connection, and \texttt{inet\_server\_port} returns the port number. All these functions return NULL if the current connection is via a Unix-domain socket.

\texttt{pg\_blocking\_pids} returns an array of the process IDs of the sessions that are blocking the server process with the specified process ID, or an empty array if there is no such server process or it is not blocked. One server process blocks another if it either holds a lock that conflicts with the blocked process’s lock request (hard block), or is waiting for a lock that would conflict with the blocked process’s lock request and is ahead of it in the wait queue (soft block). When using parallel queries the result always lists client-visible process IDs (that is, \texttt{pg\_backend\_pid} results) even if the actual lock is held or awaited by a child worker process. As a result of that, there may be duplicated PIDs in the result. Also note that when a prepared transaction holds a conflicting lock, it will be represented by a zero process ID in the result of this function. Frequent calls to this function could have some impact on database performance, because it needs exclusive access to the lock manager’s shared state for a short time.

\texttt{pg\_conf\_load\_time} returns the timestamp with time zone when the server configuration files were last loaded. (If the current session was alive at the time, this will be the time when the session itself re-read the configuration files, so the reading will vary a little in different sessions. Otherwise it is the time when the postmaster process re-read the configuration files.)

\texttt{pg\_my\_temp\_schema} returns the OID of the current session’s temporary schema, or zero if it has none (because it has not created any temporary tables). \texttt{pg\_is\_other\_temp\_schema} returns true if the given OID is the OID of another session’s temporary schema. (This can be useful, for example, to exclude other sessions’ temporary tables from a catalog display.)

\texttt{pg\_listening\_channels} returns a set of names of asynchronous notification channels that the current session is listening to. \texttt{pg\_notification\_queue\_usage} returns the fraction of the total available space for notifications currently occupied by notifications that are waiting to be processed, as a double in the range 0-1. See \texttt{LISTEN} and \texttt{NOTIFY} for more information.

\texttt{pg\_postmaster\_start\_time} returns the timestamp with time zone when the server started.

\texttt{version} returns a string describing the PostgreSQL server’s version. You can also get this information from \texttt{server\_version} or for a machine-readable version, \texttt{server\_version\_num}. Software developers should use \texttt{server\_version\_num} (available since 8.2) or \texttt{PQserverVersion} instead of parsing the text version.

Table 9-60 lists functions that allow the user to query object access privileges programmatically. See Section 5.6 for more information about privileges.

\textbf{Table 9-60. Access Privilege Inquiry Functions}

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{inet_client_addr}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{inet_client_port}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{inet_server_addr}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{inet_server_port}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{pg_blocking_pids}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{pg_conf_load_time}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{pg_my_temp_schema}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{pg_is_other_temp_schema}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{pg_listening_channels}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{pg_notification_queue_usage}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{pg_postmaster_start_time}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>\texttt{version}</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>------------------------------------------</td>
<td>-------------</td>
<td>--------------------------------------------------</td>
</tr>
<tr>
<td>has_any_column_privilege(user, table, privilege)</td>
<td>boolean</td>
<td>does user have privilege for any column of table</td>
</tr>
<tr>
<td>has_any_column_privilege(table, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for any column of table</td>
</tr>
<tr>
<td>has_column_privilege(user, table, column, privilege)</td>
<td>boolean</td>
<td>does user have privilege for column</td>
</tr>
<tr>
<td>has_column_privilege(table, column, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for column</td>
</tr>
<tr>
<td>has_database_privilege(user, database, privilege)</td>
<td>boolean</td>
<td>does user have privilege for database</td>
</tr>
<tr>
<td>has_database_privilege(database, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for database</td>
</tr>
<tr>
<td>has_foreign_data_wrapper_privilege(user, fdw, privilege)</td>
<td>boolean</td>
<td>does user have privilege for foreign-data wrapper</td>
</tr>
<tr>
<td>has_foreign_data_wrapper_privilege(fdw, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for foreign-data wrapper</td>
</tr>
<tr>
<td>has_function_privilege(user, function, privilege)</td>
<td>boolean</td>
<td>does user have privilege for function</td>
</tr>
<tr>
<td>has_function_privilege(function, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for function</td>
</tr>
<tr>
<td>has_language_privilege(user, language, privilege)</td>
<td>boolean</td>
<td>does user have privilege for language</td>
</tr>
<tr>
<td>has_language_privilege(language, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for language</td>
</tr>
<tr>
<td>has_schema_privilege(user, schema, privilege)</td>
<td>boolean</td>
<td>does user have privilege for schema</td>
</tr>
<tr>
<td>has_schema_privilege(schema, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for schema</td>
</tr>
<tr>
<td>has_sequence_privilege(user, sequence, privilege)</td>
<td>boolean</td>
<td>does user have privilege for sequence</td>
</tr>
<tr>
<td>has_sequence_privilege(sequence, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for sequence</td>
</tr>
<tr>
<td>has_server_privilege(user, server, privilege)</td>
<td>boolean</td>
<td>does user have privilege for foreign server</td>
</tr>
<tr>
<td>has_server_privilege(server, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for foreign server</td>
</tr>
<tr>
<td>has_table_privilege(user, table, privilege)</td>
<td>boolean</td>
<td>does user have privilege for table</td>
</tr>
<tr>
<td>has_table_privilege(table, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for table</td>
</tr>
<tr>
<td>has_tablespace_privilege(user, tablespace, privilege)</td>
<td>boolean</td>
<td>does user have privilege for tablespace</td>
</tr>
<tr>
<td>has_tablespace_privilege(tablespace, privilege)</td>
<td>boolean</td>
<td>does current user have privilege for tablespace</td>
</tr>
</tbody>
</table>
has_table_privilege checks whether a user can access a table in a particular way. The user can be specified by name, by OID (pg_authid.oid), public to indicate the PUBLIC pseudo-role, or if the argument is omitted current_user is assumed. The table can be specified by name or by OID. (Thus, there are actually six variants of has_table_privilege, which can be distinguished by the number and types of their arguments.) When specifying by name, the name can be schema-qualified if necessary. The desired access privilege type is specified by a text string, which must evaluate to one of the values SELECT, INSERT, UPDATE, DELETE, TRUNCATE, REFERENCES, or TRIGGER. Optionally, WITH GRANT OPTION can be added to a privilege type to test whether the privilege is held with grant option. Also, multiple privilege types can be listed separated by commas, in which case the result will be true if any of the listed privileges is held. (Case of the privilege string is not significant, and extra whitespace is allowed between but not within privilege names.) Some examples:

```sql
SELECT has_table_privilege('myschema.mytable', 'select');
SELECT has_table_privilege('joe', 'mytable', 'INSERT, SELECT WITH GRANT OPTION');
```

has_sequence_privilege checks whether a user can access a sequence in a particular way. The possibilities for its arguments are analogous to has_table_privilege. The desired access privilege type must evaluate to one of USAGE, SELECT, or UPDATE.

has_any_column_privilege checks whether a user can access any column of a table in a particular way. Its argument possibilities are analogous to has_table_privilege, except that the desired access privilege type must evaluate to some combination of SELECT, INSERT, UPDATE, or REFERENCES. Note that having any of these privileges at the table level implicitly grants it for each column of the table, so has_any_column_privilege will always return true if has_table_privilege does for the same arguments. But has_any_column_privilege also succeeds if there is a column-level grant of the privilege for at least one column.

has_column_privilege checks whether a user can access a column in a particular way. Its argument possibilities are analogous to has_table_privilege, with the addition that the column can be specified either by name or attribute number. The desired access privilege type must evaluate to some combination of SELECT, INSERT, UPDATE, or REFERENCES. Note that having any of these privileges at the table level implicitly grants it for each column of the table.

has_database_privilege checks whether a user can access a database in a particular way. Its argument possibilities are analogous to has_table_privilege. The desired access privilege type must evaluate to some combination of CREATE, CONNECT, TEMPORARY, or TEMP (which is equivalent to TEMPORARY).

has_function_privilege checks whether a user can access a function in a particular way. Its argument possibilities are analogous to has_table_privilege. When specifying a function by a
text string rather than by OID, the allowed input is the same as for the regprocedure data type (see Section 8.18). The desired access privilege type must evaluate to EXECUTE. An example is:

```sql
SELECT has_function_privilege('joeuser', 'myfunc(int, text)', 'execute');
```

`has_foreign_data_wrapper_privilege` checks whether a user can access a foreign-data wrapper in a particular way. Its argument possibilities are analogous to `has_table_privilege`. The desired access privilege type must evaluate to `USAGE`.

`has_language_privilege` checks whether a user can access a procedural language in a particular way. Its argument possibilities are analogous to `has_table_privilege`. The desired access privilege type must evaluate to `USAGE`.

`has_schema_privilege` checks whether a user can access a schema in a particular way. Its argument possibilities are analogous to `has_table_privilege`. The desired access privilege type must evaluate to some combination of `CREATE` or `USAGE`.

`has_server_privilege` checks whether a user can access a foreign server in a particular way. Its argument possibilities are analogous to `has_table_privilege`. The desired access privilege type must evaluate to `USAGE`.

`has_tablespace_privilege` checks whether a user can access a tablespace in a particular way. Its argument possibilities are analogous to `has_table_privilege`. The desired access privilege type must evaluate to `CREATE`.

`has_type_privilege` checks whether a user can access a type in a particular way. Its argument possibilities are analogous to `has_table_privilege`, except that `public` is not allowed as a user name. The desired access privilege type must evaluate to some combination of `MEMBER` or `USAGE`. `MEMBER` denotes direct or indirect membership in the role (that is, the right to do `SET ROLE`), while `USAGE` denotes whether the privileges of the role are immediately available without doing `SET ROLE`.

`row_security_active` checks whether row level security is active for the specified table in the context of the `current_user` and environment. The table can be specified by name or by OID.

Table 9-61 shows functions that determine whether a certain object is visible in the current schema search path. For example, a table is said to be visible if its containing schema is in the search path and no table of the same name appears earlier in the search path. This is equivalent to the statement that the table can be referenced by name without explicit schema qualification. To list the names of all visible tables:

```sql
SELECT relname FROM pg_class WHERE pg_table_is_visible(oid);
```

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_collation_is_visible</code></td>
<td>boolean</td>
<td>is collation visible in search path</td>
</tr>
<tr>
<td><code>pg_conversion_is_visible</code></td>
<td>boolean</td>
<td>is conversion visible in search path</td>
</tr>
</tbody>
</table>

Table 9-61. Schema Visibility Inquiry Functions
Each function performs the visibility check for one type of database object. Note that `pg_table_is_visible` can also be used with views, materialized views, indexes, sequences and foreign tables; `pg_type_is_visible` can also be used with domains. For functions and operators, an object in the search path is visible if there is no object of the same name and argument data type(s) earlier in the path. For operator classes, both name and associated index access method are considered.

All these functions require object OIDs to identify the object to be checked. If you want to test an object by name, it is convenient to use the OID alias types (regclass, regtype, regprocedure, regoperator, regconfig, or regdictionary), for example:

```sql
SELECT pg_type_is_visible('myschema.widget'::regtype);
```

Note that it would not make much sense to test a non-schema-qualified type name in this way — if the name can be recognized at all, it must be visible.

Table 9-62 lists functions that extract information from the system catalogs.

## Table 9-62. System Catalog Information Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>format_type(type_oid, typemod)</td>
<td>text</td>
<td>get SQL name of a data type</td>
</tr>
<tr>
<td>pg_get_constraintdef(constraint_oid)</td>
<td>text</td>
<td>get definition of a constraint</td>
</tr>
<tr>
<td>pg_get_constraintdef(constraint_oid, pretty_bool)</td>
<td>text, bool</td>
<td>get definition of a constraint</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------------------</td>
<td>-------------</td>
<td>------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>pg_get_expr(pg_node_tree, relation_oid)</td>
<td>text</td>
<td>decompile internal form of an expression, assuming that any Vars in it refer to the relation indicated by the second parameter</td>
</tr>
<tr>
<td>pg_get_expr(pg_node_tree, relation_oid, pretty_bool)</td>
<td>text</td>
<td>decompile internal form of an expression, assuming that any Vars in it refer to the relation indicated by the second parameter</td>
</tr>
<tr>
<td>pg_get_functiondef(func_oid)</td>
<td>text</td>
<td>get definition of a function</td>
</tr>
<tr>
<td>pg_get_function_arguments(func_oid)</td>
<td>text</td>
<td>get argument list of function’s definition (with default values)</td>
</tr>
<tr>
<td>pg_get_function_identity_arguments(func_oid)</td>
<td>text</td>
<td>get argument list to identify a function (without default values)</td>
</tr>
<tr>
<td>pg_get_function_result(func_oid)</td>
<td>text</td>
<td>get RETURNS clause for function</td>
</tr>
<tr>
<td>pg_get_indexdef(index_oid)</td>
<td>text</td>
<td>get CREATE INDEX command for index</td>
</tr>
<tr>
<td>pg_get_indexdef(index_oid, column_no, pretty_bool)</td>
<td>text</td>
<td>get CREATE INDEX command for index, or definition of just one index column when column_no is not zero</td>
</tr>
<tr>
<td>pg_get_keywords()</td>
<td>setof record</td>
<td>get list of SQL keywords and their categories</td>
</tr>
<tr>
<td>pg_get_ruledef(rule_oid)</td>
<td>text</td>
<td>get CREATE RULE command for rule</td>
</tr>
<tr>
<td>pg_get_ruledef(rule_oid, pretty_bool)</td>
<td>text</td>
<td>get CREATE RULE command for rule</td>
</tr>
<tr>
<td>pg_get_serial_sequence(table_name, column_name)</td>
<td>text</td>
<td>get name of the sequence that a serial, smallserial or bigserial column uses</td>
</tr>
<tr>
<td>pg_get_triggerdef(trigger_name)</td>
<td>text</td>
<td>get CREATE [ CONSTRAINT ] TRIGGER command for trigger</td>
</tr>
<tr>
<td>pg_get_triggerdef(trigger_name, pretty_bool)</td>
<td>text</td>
<td>get CREATE [ CONSTRAINT ] TRIGGER command for trigger</td>
</tr>
<tr>
<td>pg_get_userbyid(role_oid)</td>
<td>name</td>
<td>get role name with given OID</td>
</tr>
<tr>
<td>pg_get_viewdef(view_name)</td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view (deprecated)</td>
</tr>
<tr>
<td>pg_get_viewdef(view_name, pretty_bool)</td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view (deprecated)</td>
</tr>
</tbody>
</table>
### Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_get_viewdef(view_oid)</code></td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view</td>
</tr>
<tr>
<td><code>pg_get_viewdef(view_oid, pretty_bool)</code></td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view</td>
</tr>
<tr>
<td><code>pg_get_viewdef(view_oid, wrap_column_int)</code></td>
<td>text</td>
<td>get underlying SELECT command for view or materialized view; lines with</td>
</tr>
<tr>
<td></td>
<td></td>
<td>fields are wrapped to specified number of columns, pretty-printing is implied</td>
</tr>
<tr>
<td><code>pg_index_column_has_property(index_oid, column_no, prop_name)</code></td>
<td>boolean</td>
<td>test whether an index column has a specified property</td>
</tr>
<tr>
<td><code>pg_index_has_property(index_oid, prop_name)</code></td>
<td>boolean</td>
<td>test whether an index has a specified property</td>
</tr>
<tr>
<td><code>pg_indexam_has_property(am_oid, prop_name)</code></td>
<td>boolean</td>
<td>test whether an index access method has a specified property</td>
</tr>
<tr>
<td><code>pg_options_to_table(reloptions)</code></td>
<td>setof record</td>
<td>get the set of storage option name/value pairs</td>
</tr>
<tr>
<td><code>pg_tablespace_databases(tablespace_oid)</code></td>
<td>setof oid</td>
<td>get the set of database OIDs that have objects in the tablespace</td>
</tr>
<tr>
<td><code>pg_tablespace_location(tablespace_oid)</code></td>
<td>text</td>
<td>get the path in the file system that this tablespace is located in</td>
</tr>
<tr>
<td><code>pg_typeof(any)</code></td>
<td>regtype</td>
<td>get the data type of any value</td>
</tr>
<tr>
<td><code>collation_for(any)</code></td>
<td>text</td>
<td>get the collation of the argument</td>
</tr>
<tr>
<td><code>to_regclass(rel_name)</code></td>
<td>regclass</td>
<td>get the OID of the named relation</td>
</tr>
<tr>
<td><code>to_regproc(func_name)</code></td>
<td>regproc</td>
<td>get the OID of the named function</td>
</tr>
<tr>
<td><code>to_regprocedure(func_name)</code></td>
<td>regprocedure</td>
<td>get the OID of the named function</td>
</tr>
<tr>
<td><code>to_regoper(operator_name)</code></td>
<td>regoper</td>
<td>get the OID of the named operator</td>
</tr>
<tr>
<td><code>to_regoperator(operator_name)</code></td>
<td>regoperator</td>
<td>get the OID of the named operator</td>
</tr>
<tr>
<td><code>to_regtype(type_name)</code></td>
<td>regtype</td>
<td>get the OID of the named type</td>
</tr>
<tr>
<td><code>to_regnamespace(schema_name)</code></td>
<td>regnamespace</td>
<td>get the OID of the named schema</td>
</tr>
<tr>
<td><code>to_regrole(role_name)</code></td>
<td>regrole</td>
<td>get the OID of the named role</td>
</tr>
</tbody>
</table>

`format_type` returns the SQL name of a data type that is identified by its type OID and possibly a type modifier. Pass NULL for the type modifier if no specific modifier is known.
pg_get_keywords returns a set of records describing the SQL keywords recognized by the server. The word column contains the keyword. The catcode column contains a category code: U for unreserved, C for column name, T for type or function name, or R for reserved. The catdesc column contains a possibly-localized string describing the category.

pg_get_constraintdef, pg_get_indexdef, pg_get_ruledef, and pg_get_triggerdef, respectively reconstruct the creating command for a constraint, index, rule, or trigger. (Note that this is a decompiled reconstruction, not the original text of the command.) pg_get_expr decompiles the internal form of an individual expression, such as the default value for a column. It can be useful when examining the contents of system catalogs. If the expression might contain Vars, specify the OID of the relation they refer to as the second parameter; if no Vars are expected, zero is sufficient. pg_get_viewdef reconstructs the SELECT query that defines a view. Most of these functions come in two variants, one of which can optionally “pretty-print” the result. The pretty-printed format is more readable, but the default format is more likely to be interpreted the same way by future versions of PostgreSQL; avoid using pretty-printed output for dump purposes. Passing false for the pretty-print parameter yields the same result as the variant that does not have the parameter at all.

pg_get_functiondef returns a complete CREATE OR REPLACE FUNCTION statement for a function. pg_get_function_arguments returns the argument list of a function, in the form it would need to appear in within CREATE FUNCTION. pg_get_function_result similarly returns the appropriate RETURNS clause for the function. pg_get_function_identity_arguments returns the argument list necessary to identify a function, in the form it would need to appear in within ALTER FUNCTION, for instance. This form omits default values.

pg_get_serial_sequence returns the name of the sequence associated with a column, or NULL if no sequence is associated with the column. The first input parameter is a table name with optional schema, and the second parameter is a column name. Because the first parameter is potentially a schema and table, it is not treated as a double-quoted identifier, meaning it is lower cased by default, while the second parameter, being just a column name, is treated as double-quoted and has its case preserved. The function returns a value suitably formatted for passing to sequence functions (see Section 9.16). This association can be modified or removed with ALTER SEQUENCE OWNED BY. (The function probably should have been called pg_get_owned_sequence; its current name reflects the fact that it’s typically used with serial or bigserial columns.)

pg_get_userbyid extracts a role’s name given its OID.

pg_index_column_has_property, pg_index_has_property, and pg_indexam_has_property return whether the specified index column, index, or index access method possesses the named property. NULL is returned if the property name is not known or does not apply to the particular object, or if the OID or column number does not identify a valid object. Refer to Table 9-63 for column properties, Table 9-64 for index properties, and Table 9-65 for access method properties. (Note that extension access methods can define additional property names for their indexes.)

Table 9-63. Index Column Properties

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>asc</td>
<td>Does the column sort in ascending order on a forward scan?</td>
</tr>
<tr>
<td>desc</td>
<td>Does the column sort in descending order on a forward scan?</td>
</tr>
<tr>
<td>nulls_first</td>
<td>Does the column sort with nulls first on a forward scan?</td>
</tr>
<tr>
<td>Name</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>nulls_last</td>
<td>Does the column sort with nulls last on a forward scan?</td>
</tr>
<tr>
<td>orderable</td>
<td>Does the column possess any defined sort ordering?</td>
</tr>
<tr>
<td>distance_orderable</td>
<td>Can the column be scanned in order by a “distance” operator, for example ORDER BY col &lt;= constant?</td>
</tr>
<tr>
<td>returnable</td>
<td>Can the column value be returned by an index-only scan?</td>
</tr>
<tr>
<td>search_array</td>
<td>Does the column natively support col = ANY(array) searches?</td>
</tr>
<tr>
<td>search_nulls</td>
<td>Does the column support IS NULL and IS NOT NULL searches?</td>
</tr>
</tbody>
</table>

Table 9-64. Index Properties

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>clusterable</td>
<td>Can the index be used in a CLUSTER command?</td>
</tr>
<tr>
<td>index_scan</td>
<td>Does the index support plain (non-bitmap) scans?</td>
</tr>
<tr>
<td>bitmap_scan</td>
<td>Does the index support bitmap scans?</td>
</tr>
<tr>
<td>backward_scan</td>
<td>Can the scan direction be changed in mid-scan (to support FETCH BACKWARD on a cursor without needing materialization)?</td>
</tr>
</tbody>
</table>

Table 9-65. Index Access Method Properties

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>can_order</td>
<td>Does the access method support ASC, DESC and related keywords in CREATE_INDEX?</td>
</tr>
<tr>
<td>can_unique</td>
<td>Does the access method support unique indexes?</td>
</tr>
<tr>
<td>can_multi_col</td>
<td>Does the access method support indexes with multiple columns?</td>
</tr>
<tr>
<td>can_exclude</td>
<td>Does the access method support exclusion constraints?</td>
</tr>
</tbody>
</table>

pg_options_to_table returns the set of storage option name/value pairs (option_name/option_value) when passed pg_class.reloptions or pg_attribute.attoptions.

pg_tablespace_databases allows a tablespace to be examined. It returns the set of OIDs of databases that have objects stored in the tablespace. If this function returns any rows, the tablespace is not empty and cannot be dropped. To display the specific objects populating the tablespace, you will need to connect to the databases identified by pg_tablespace_databases and query their pg_class catalogs.

pg_typeof returns the OID of the data type of the value that is passed to it. This can be helpful
for troubleshooting or dynamically constructing SQL queries. The function is declared as returning a type named `regtype`, which is an OID alias type (see Section 8.18); this means that it is the same as an OID for comparison purposes but displays as a type name. For example:

```sql
SELECT pg_typeof(33);
```

```
pg_typeof
---------
integer
(1 row)
```

```sql
SELECT typlen FROM pg_type WHERE oid = pg_typeof(33);
```

```
typlen
--------
     4
(1 row)
```

The expression `collation for` returns the collation of the value that is passed to it. Example:

```sql
SELECT collation for (description) FROM pg_description LIMIT 1;
```

```
p_collation_for
---------------
"default"
(1 row)
```

```sql
SELECT collation for ('foo' COLLATE "de_DE");
```

```
p_collation_for
---------------
"de_DE"
(1 row)
```

The value might be quoted and schema-qualified. If no collation is derived for the argument expression, then a null value is returned. If the argument is not of a collatable data type, then an error is raised.

The `to_regclass`, `to_regproc`, `to_regprocedure`, `to_regoper`, `to_regoperator`, `to_regtype`, `to_regnamespace`, and `to_regrole` functions translate relation, function, operator, type, schema, and role names (given as `text`) to objects of type `regclass`, `regproc`, `regprocedure`, `regoper`, `regoperator`, `regtype`, `regnamespace`, and `regrole` respectively. These functions differ from a cast from text in that they don’t accept a numeric OID, and that they return null rather than throwing an error if the name is not found (or, for `to_regproc` and `to_regoper`, if the given name matches multiple objects).

Table 9-66 lists functions related to database object identification and addressing.

### Table 9-66. Object Information and Addressing Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_describe_object(classid oid, objid oid, objsubid integer)</code></td>
<td><code>text</code></td>
<td>get description of a database object</td>
</tr>
<tr>
<td><code>pg_identify_object(classid oid, objid oid, objsubid integer)</code></td>
<td><code>type text, schema text, name text, identity text</code></td>
<td>get identity of a database object</td>
</tr>
</tbody>
</table>
pg_identify_object_as_address returns a row containing enough information to uniquely identify the database object specified by catalog OID, object OID, and sub-object ID (such as a column number within a table; the sub-object ID is zero when referring to a whole object). This description is intended to be human-readable, and might be translated, depending on server configuration. This is useful to determine the identity of an object as stored in the pg_depend catalog.

pg_identify_object returns a row containing enough information to uniquely identify the database object specified by catalog OID, object OID and sub-object ID. This information is intended to be machine-readable, and is never translated. type identifies the type of database object; schema is the schema name that the object belongs in, or NULL for object types that do not belong to schemas; name is the name of the object, quoted if necessary, if the name (along with schema name, if pertinent) is sufficient to uniquely identify the object, otherwise NULL; identity is the complete object identity, with the precise format depending on object type, and each name within the format being schema-qualified and quoted as necessary.

pg_identify_object_as_address returns a row containing enough information to uniquely identify the database object specified by catalog OID, object OID and sub-object ID. The returned information is independent of the current server, that is, it could be used to identify an identically named object in another server. type identifies the type of database object; object_names and object_args are text arrays that together form a reference to the object. These three values can be passed to pg_get_object_address to obtain the internal address of the object. This function is the inverse of pg_identify_object_as_address.

pg_get_object_address returns a row containing enough information to uniquely identify the database object specified by its type and object name and argument arrays. The returned values are the ones that would be used in system catalogs such as pg_depend and can be passed to other system functions such as pg_identify_object or pg_describe_object. classid is the OID of the system catalog containing the object; objid is the OID of the object itself, and objsubid is the sub-object ID, or zero if none. This function is the inverse of pg_identify_object_as_address.

The functions shown in Table 9-67 extract comments previously stored with the COMMENT command. A null value is returned if no comment could be found for the specified parameters.

### Table 9-67. Comment Information Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>col_description(table_oid, column_number)</td>
<td>text</td>
<td>get comment for a table column</td>
</tr>
<tr>
<td>obj_description(object_oid, catalog_name)</td>
<td>text</td>
<td>get comment for a database object</td>
</tr>
<tr>
<td>obj_description(object_oid)</td>
<td>text</td>
<td>get comment for a database object   (deprecated)</td>
</tr>
<tr>
<td>shobj_description(object_oid, catalog_name)</td>
<td>text</td>
<td>get comment for a shared database object</td>
</tr>
</tbody>
</table>
col_description returns the comment for a table column, which is specified by the OID of its table and its column number. (obj_description cannot be used for table columns since columns do not have OIDs of their own.)

The two-parameter form of obj_description returns the comment for a database object specified by its OID and the name of the containing system catalog. For example, obj_description(123456, 'pg_class') would retrieve the comment for the table with OID 123456. The one-parameter form of obj_description requires only the object OID. It is deprecated since there is no guarantee that OIDs are unique across different system catalogs; therefore, the wrong comment might be returned.

shobj_description is used just like obj_description except it is used for retrieving comments on shared objects. Some system catalogs are global to all databases within each cluster, and the descriptions for objects in them are stored globally as well.

The functions shown in Table 9-68 provide server transaction information in an exportable form. The main use of these functions is to determine which transactions were committed between two snapshots.

Table 9-68. Transaction IDs and Snapshots

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>txid_current()</td>
<td>bigint</td>
<td>get current transaction ID, assigning a new one if the current transaction does not have one</td>
</tr>
<tr>
<td>txid_current_snapshot()</td>
<td>txid_snapshot</td>
<td>get current snapshot</td>
</tr>
<tr>
<td>txid_snapshot_xip(txid_snapshot)</td>
<td>setof bigint</td>
<td>get in-progress transaction IDs in snapshot</td>
</tr>
<tr>
<td>txid_snapshot_xmax(txid_snapshot)</td>
<td>bigint</td>
<td>get xmax of snapshot</td>
</tr>
<tr>
<td>txid_snapshot_xmin(txid_snapshot)</td>
<td>bigint</td>
<td>get xmin of snapshot</td>
</tr>
<tr>
<td>txid_visible_in_snapshot(bigint, txid_snapshot)</td>
<td>boolean</td>
<td>is transaction ID visible in snapshot? (do not use with subtransaction ids)</td>
</tr>
</tbody>
</table>

The internal transaction ID type (xid) is 32 bits wide and wraps around every 4 billion transactions. However, these functions export a 64-bit format that is extended with an “epoch” counter so it will not wrap around during the life of an installation. The data type used by these functions, txid_snapshot, stores information about transaction ID visibility at a particular moment in time. Its components are described in Table 9-69.

Table 9-69. Snapshot Components

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>xmin</td>
<td>Earliest transaction ID (txid) that is still active. All earlier transactions will either be committed and visible, or rolled back and dead.</td>
</tr>
<tr>
<td>xmax</td>
<td>First as-yet-unassigned txid. All txids greater than or equal to this are not yet started as of the time of the snapshot, and thus invisible.</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>xip_list</td>
<td>Active txids at the time of the snapshot. The list includes only those active txids between xmin and xmax; there might be active txids higher than xmax. A txid that is xmin &lt;= txid &lt; xmax and not in this list was already completed at the time of the snapshot, and thus either visible or dead according to its commit status. The list does not include txids of subtransactions.</td>
</tr>
</tbody>
</table>

txid_snapshot’s textual representation is xmin:xmax:xip_list. For example 10:20:10,14,15 means xmin=10, xmax=20, xip_list=10, 14, 15.

The functions shown in Table 9-70 provide information about transactions that have been already committed. These functions mainly provide information about when the transactions were committed. They only provide useful data when track_commit_timestamp configuration option is enabled and only for transactions that were committed after it was enabled.

### Table 9-70. Committed transaction information

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_xact_commit_timestamp(xid)</td>
<td>timestamp with time zone</td>
<td>get commit timestamp of a transaction</td>
</tr>
<tr>
<td>pg_last_committed_xact()</td>
<td>xid xid,timestamp</td>
<td>get transaction ID and commit timestamp of latest committed transaction</td>
</tr>
</tbody>
</table>

The functions shown in Table 9-71 print information initialized during initdb, such as the catalog version. They also show information about write-ahead logging and checkpoint processing. This information is cluster-wide, and not specific to any one database. They provide most of the same information, from the same source, as pg_controldata, although in a form better suited to SQL functions.

### Table 9-71. Control Data Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_control_checkpoint()</td>
<td>record</td>
<td>Returns information about current checkpoint state.</td>
</tr>
<tr>
<td>pg_control_system()</td>
<td>record</td>
<td>Returns information about current control file state.</td>
</tr>
<tr>
<td>pg_control_init()</td>
<td>record</td>
<td>Returns information about cluster initialization state.</td>
</tr>
<tr>
<td>pg_control_recovery()</td>
<td>record</td>
<td>Returns information about recovery state.</td>
</tr>
</tbody>
</table>

*pg_control_checkpoint* returns a record, shown in Table 9-72

### Table 9-72. pg_control_checkpoint Columns

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>checkpoint_location</td>
<td>pg_lsn</td>
</tr>
</tbody>
</table>
### Column Name | Data Type
---|---
prior_location | pg_lsn
redo_location | pg_lsn
redo_wal_file | text
timeline_id | integer
prev_timeline_id | integer
full_page_writes | boolean
next_xid | text
next_oid | oid
next_multixact_id | xid
next_multi_offset | xid
oldest_xid | xid
oldest_xid_dbid | oid
oldest_active_xid | xid
oldest_multi_xid | xid
oldest_multi_dbid | oid
oldest_commit_ts_xid | xid
newest_commit_ts_xid | xid
checkpoint_time | timestamp with time zone

pg_control_system returns a record, shown in Table 9-73

#### Table 9-73. pg_control_system Columns

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
</table>
pq_control_version         | integer                |
catalog_version_no         | integer                |
system_identifier          | bigint                 |
pq_control_last_modified   | timestamp with time zone |

pg_control_init returns a record, shown in Table 9-74

#### Table 9-74. pg_control_init Columns

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
</table>
max_data_alignment          | integer              |
database_block_size         | integer              |
blocks_per_segment          | integer              |
wal_block_size              | integer              |
bytes_per_wal_segment       | integer              |
max_identifier_length      | integer              |
max_index_columns           | integer              |
max_toast_chunk_size        | integer              |
large_object_chunk_size     | integer              |
bigint_timestamps           | boolean              |
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>float4_pass_by_value</td>
<td>boolean</td>
</tr>
<tr>
<td>float8_pass_by_value</td>
<td>boolean</td>
</tr>
<tr>
<td>data_page_checksum_version</td>
<td>integer</td>
</tr>
</tbody>
</table>

**pg_control_recovery** returns a record, shown in Table 9-75.

Table 9-75. pg_control_recovery Columns

<table>
<thead>
<tr>
<th>Column Name</th>
<th>Data Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>min_recovery_end_location</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>min_recovery_end_timeline</td>
<td>integer</td>
</tr>
<tr>
<td>backup_start_location</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>backup_end_location</td>
<td>pg_lsn</td>
</tr>
<tr>
<td>end_of_backup_record_required</td>
<td>boolean</td>
</tr>
</tbody>
</table>

### 9.26. System Administration Functions

The functions described in this section are used to control and monitor a PostgreSQL installation.

#### 9.26.1. Configuration Settings Functions

Table 9-76 shows the functions available to query and alter run-time configuration parameters.

Table 9-76. Configuration Settings Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>current_setting(setting_name [, missing_ok])</td>
<td>text</td>
<td>get current value of setting</td>
</tr>
<tr>
<td>set_config(setting_name, new_value, is_local)</td>
<td>text</td>
<td>set parameter and return new value</td>
</tr>
</tbody>
</table>

The function `current_setting` yields the current value of the setting `setting_name`. It corresponds to the SQL command `SHOW`. An example:

```sql
SELECT current_setting('datestyle');
```

```plaintext
current_setting
-------------
ISO, MDY
(1 row)
```

If there is no setting named `setting_name`, `current_setting` throws an error unless `missing_ok` is supplied and is `true`.

`set_config` sets the parameter `setting_name` to `new_value`. If `is_local` is `true`, the `new_value` will only apply to the current transaction. If you want the new value to apply for the current session, use `false` instead. The function corresponds to the SQL command `SET`. An example:

```sql
set_config(setting_name, new_value, is_local)
```

```plaintext
set_config
----------
```

If there is no `setting_name`, `set_config` throws an error unless `is_local` is `true`.
9.26.2. Server Signaling Functions

The functions shown in Table 9-77 send control signals to other server processes. Use of these functions is restricted to superusers by default but access may be granted to others with the GRANT, with noted exceptions.

Table 9-77. Server Signaling Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_cancel_backend(pid int)</td>
<td>boolean</td>
<td>Cancel a backend’s current query. This is also allowed if the calling role is a member of the role whose backend is being canceled or the calling role has been granted pg_signal_backend, however only superusers can cancel superuser backends.</td>
</tr>
<tr>
<td>pg_reload_conf()</td>
<td>boolean</td>
<td>Cause server processes to reload their configuration files</td>
</tr>
<tr>
<td>pg_rotate_logfile()</td>
<td>boolean</td>
<td>Rotate server’s log file</td>
</tr>
<tr>
<td>pg_terminate_backend(pid int)</td>
<td>boolean</td>
<td>Terminate a backend. This is also allowed if the calling role is a member of the role whose backend is being terminated or the calling role has been granted pg_signal_backend, however only superusers can terminate superuser backends.</td>
</tr>
</tbody>
</table>

Each of these functions returns true if successful and false otherwise.

pg_cancel_backend and pg_terminate_backend send signals (SIGINT or SIGTERM respectively) to backend processes identified by process ID. The process ID of an active backend can be found from the pid column of the pg_stat_activity view, or by listing the postgres processes on the server (using ps on Unix or the Task Manager on Windows). The role of an active backend can be found from the usename column of the pg_stat_activity view.

pg_reload_conf sends a SIGHUP signal to the server, causing configuration files to be reloaded by all server processes.

pg_rotate_logfile signals the log-file manager to switch to a new output file immediately. This works only when the built-in log collector is running, since otherwise there is no log-file manager subprocess.
9.26.3. Backup Control Functions

The functions shown in Table 9-78 assist in making on-line backups. These functions cannot be executed during recovery (except non-exclusive `pg_start_backup`, non-exclusive `pg_stop_backup`, `pg_is_in_backup`, `pg_backup_start_time` and `pg_xlog_location_diff`).

Table 9-78. Backup Control Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_create_restore_point(name text)</code></td>
<td><code>pg_lsn</code></td>
<td>Create a named point for performing restore (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td><code>pg_current_xlog_flush_location()</code></td>
<td><code>pg_lsn</code></td>
<td>Get current transaction log flush location</td>
</tr>
<tr>
<td><code>pg_current_xlog_insert_location()</code></td>
<td><code>pg_lsn</code></td>
<td>Get current transaction log insert location</td>
</tr>
<tr>
<td><code>pg_current_xlog_location()</code></td>
<td><code>pg_lsn</code></td>
<td>Get current transaction log write location</td>
</tr>
<tr>
<td><code>pg_start_backup(label text [, fast boolean [, exclusive boolean ]])</code></td>
<td><code>pg_lsn</code></td>
<td>Prepare for performing on-line backup (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td><code>pg_stop_backup()</code></td>
<td><code>pg_lsn</code></td>
<td>Finish performing exclusive on-line backup (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td><code>pg_stop_backup(exclusive boolean)</code></td>
<td><code>setof record</code></td>
<td>Finish performing exclusive or non-exclusive on-line backup (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
<tr>
<td><code>pg_is_in_backup()</code></td>
<td><code>bool</code></td>
<td>True if an on-line exclusive backup is still in progress.</td>
</tr>
<tr>
<td><code>pg_backup_start_time()</code></td>
<td><code>timestamp with time zone</code></td>
<td>Get start time of an on-line exclusive backup in progress.</td>
</tr>
<tr>
<td><code>pg_switch_xlog()</code></td>
<td><code>pg_lsn</code></td>
<td>Force switch to a new transaction log file (restricted to superusers by default, but other users can be granted EXECUTE to run the function)</td>
</tr>
</tbody>
</table>
Chapter 9. Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_xlogfile_name(location pg_lsn)</td>
<td>text</td>
<td>Convert transaction log location string to file name</td>
</tr>
<tr>
<td>pg_xlogfile_name_offset(location pg_lsn)</td>
<td>text, integer</td>
<td>Convert transaction log location string to file name and decimal byte offset within file</td>
</tr>
<tr>
<td>pg_xlog_location_diff(location pg_lsn, location pg_lsn)</td>
<td>numeric</td>
<td>Calculate the difference between two transaction log locations</td>
</tr>
</tbody>
</table>

pg_start_backup accepts an arbitrary user-defined label for the backup. (Typically this would be the name under which the backup dump file will be stored.) When used in exclusive mode, the function writes a backup label file (backup_label) and, if there are any links in the pg_tblspc/ directory, a tablespace map file (tablespace_map) into the database cluster’s data directory, performs a checkpoint, and then returns the backup’s starting transaction log location as text. The user can ignore this result value, but it is provided in case it is useful. When used in non-exclusive mode, the contents of these files are instead returned by the pg_stop_backup function, and should be written to the backup by the caller.

```sql
postgres=# select pg_start_backup('label_goes here');
pag_start_backup-----------------
|                      0/D4445B8 |
(1 row)
```

There is an optional second parameter of type boolean. If true, it specifies executing pg_start_backup as quickly as possible. This forces an immediate checkpoint which will cause a spike in I/O operations, slowing any concurrently executing queries.

In an exclusive backup, pg_stop_backup removes the label file and, if it exists, the tablespace_map file created by pg_start_backup. In a non-exclusive backup, the contents of the backup_label and tablespace_map are returned in the result of the function, and should be written to files in the backup (and not in the data directory). When executed on a primary pg_stop_backup will wait for WAL to be archived, provided that archiving is enabled.

On a standby pg_stop_backup will return immediately without waiting, so it’s important to verify that all required WAL segments have been archived. If write activity on the primary is low, it may be useful to run pg_switch_xlog on the primary in order to trigger a segment switch.

When executed on a primary, the function also creates a backup history file in the write-ahead log archive area. The history file includes the label given to pg_start_backup, the starting and ending transaction log locations for the backup, and the starting and ending times of the backup. The return value is the backup’s ending transaction log location (which again can be ignored). After recording the ending location, the current transaction log insertion point is automatically advanced to the next transaction log file, so that the ending transaction log file can be archived immediately to complete the backup.

pg_switch_xlog moves to the next transaction log file, allowing the current file to be archived (assuming you are using continuous archiving). The return value is the ending transaction log location + 1 within the just-completed transaction log file. If there has been no transaction log activity since the last transaction log switch, pg_switch_xlog does nothing and returns the start location of the transaction log file currently in use.

pg_create_restore_point creates a named transaction log record that can be used as recovery target, and returns the corresponding transaction log location. The given name can then be used with
recovery_target_name to specify the point up to which recovery will proceed. Avoid creating multiple restore points with the same name, since recovery will stop at the first one whose name matches the recovery target.

pg_current_xlog_location displays the current transaction log write location in the same format used by the above functions. Similarly, pg_current_xlog_insert_location displays the current transaction log insertion point and pg_current_xlog_flush_location displays the current transaction log flush point. The insertion point is the "logical" end of the transaction log at any instant, while the write location is the end of what has actually been written out from the server’s internal buffers and flush location is the location guaranteed to be written to durable storage. The write location is the end of what can be examined from outside the server, and is usually what you want if you are interested in archiving partially-complete transaction log files. The insertion and flush points are made available primarily for server debugging purposes. These are both read-only operations and do not require superuser permissions.

You can use pg_xlogfile_name_offset to extract the corresponding transaction log file name and byte offset from the results of any of the above functions. For example:

```sql
postgres=# SELECT * FROM pg_xlogfile_name_offset(pg_stop_backup());
    file_name |     file_offset
---------------+---------------
 000000100000000000000000D | 4039624
(1 row)
```

Similarly, pg_xlogfile_name extracts just the transaction log file name. When the given transaction log location is exactly at a transaction log file boundary, both these functions return the name of the preceding transaction log file. This is usually the desired behavior for managing transaction log archiving behavior, since the preceding file is the last one that currently needs to be archived.

pg_xlog_location_diff calculates the difference in bytes between two transaction log locations. It can be used with pg_stat_replication or some functions shown in Table 9-78 to get the replication lag.

For details about proper usage of these functions, see Section 25.3.

### 9.26.4. Recovery Control Functions

The functions shown in Table 9-79 provide information about the current status of the standby. These functions may be executed both during recovery and in normal running.

#### Table 9-79. Recovery Information Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_is_in_recovery()</td>
<td>bool</td>
<td>True if recovery is still in progress.</td>
</tr>
</tbody>
</table>
### Functions and Operators

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_last_xlog_receive_location()</td>
<td>pg_lsn</td>
<td>Get last transaction log location received and synced to disk by streaming replication. While streaming replication is in progress this will increase monotonically. If recovery has completed this will remain static at the value of the last WAL record received and synced to disk during recovery. If streaming replication is disabled, or if it has not yet started, the function returns NULL.</td>
</tr>
<tr>
<td>pg_last_xlog_replay_location()</td>
<td>pg_lsn</td>
<td>Get last transaction log location replayed during recovery. If recovery is still in progress this will increase monotonically. If recovery has completed then this value will remain static at the value of the last WAL record applied during that recovery. When the server has been started normally without recovery the function returns NULL.</td>
</tr>
<tr>
<td>pg_last_xact_replay_timestamp()</td>
<td>timestamp with time</td>
<td>Get time stamp of last transaction replayed during recovery. This is the time at which the commit or abort WAL record for that transaction was generated on the primary. If no transactions have been replayed during recovery, this function returns NULL. Otherwise, if recovery is still in progress this will increase monotonically. If recovery has completed then this value will remain static at the value of the last transaction applied during that recovery. When the server has been started normally without recovery the function returns NULL.</td>
</tr>
</tbody>
</table>

The functions shown in Table 9-80 control the progress of recovery. These functions may be executed only during recovery.
Table 9-80. Recovery Control Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_is_xlog_replay_paused()</td>
<td>bool</td>
<td>True if recovery is paused.</td>
</tr>
<tr>
<td>pg_xlog_replay_pause()</td>
<td>void</td>
<td>Pauses recovery immediately.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>(restricted to superusers by default, but other users can be granted EXECUTE to run the function).</td>
</tr>
<tr>
<td>pg_xlog_replay_resume()</td>
<td>void</td>
<td>Restarts recovery if it was paused (restricted to superusers by default, but other users can be granted EXECUTE to run the function).</td>
</tr>
</tbody>
</table>

While recovery is paused no further database changes are applied. If in hot standby, all new queries will see the same consistent snapshot of the database, and no further query conflicts will be generated until recovery is resumed.

If streaming replication is disabled, the paused state may continue indefinitely without problem. While streaming replication is in progress WAL records will continue to be received, which will eventually fill available disk space, depending upon the duration of the pause, the rate of WAL generation and available disk space.

9.26.5. Snapshot Synchronization Functions

PostgreSQL allows database sessions to synchronize their snapshots. A snapshot determines which data is visible to the transaction that is using the snapshot. Synchronized snapshots are necessary when two or more sessions need to see identical content in the database. If two sessions just start their transactions independently, there is always a possibility that some third transaction commits between the executions of the two START TRANSACTION commands, so that one session sees the effects of that transaction and the other does not.

To solve this problem, PostgreSQL allows a transaction to export the snapshot it is using. As long as the exporting transaction remains open, other transactions can import its snapshot, and thereby be guaranteed that they see exactly the same view of the database that the first transaction sees. But note that any database changes made by any one of these transactions remain invisible to the other transactions, as is usual for changes made by uncommitted transactions. So the transactions are synchronized with respect to pre-existing data, but act normally for changes they make themselves.

Snapshots are exported with the pg_export_snapshot function, shown in Table 9-81, and imported with the SET TRANSACTION command.

Table 9-81. Snapshot Synchronization Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_export_snapshot()</td>
<td>text</td>
<td>Save the current snapshot and return its identifier</td>
</tr>
</tbody>
</table>

The function pg_export_snapshot saves the current snapshot and returns a text string identifying
the snapshot. This string must be passed (outside the database) to clients that want to import the snapshot. The snapshot is available for import only until the end of the transaction that exported it. A transaction can export more than one snapshot, if needed. Note that doing so is only useful in READ COMMITTED transactions, since in REPEATABLE READ and higher isolation levels, transactions use the same snapshot throughout their lifetime. Once a transaction has exported any snapshots, it cannot be prepared with PREPARE TRANSACTION.

See SET TRANSACTION for details of how to use an exported snapshot.

### 9.26.6. Replication Functions

The functions shown in Table 9-82 are for controlling and interacting with replication features. See Section 26.2.5, Section 26.2.6, and Chapter 48 for information about the underlying features. Use of functions for replication origin is restricted to superusers. Use of functions for replication slot is restricted to superusers and users having REPLICATION privilege.

Many of these functions have equivalent commands in the replication protocol; see Section 51.3.

The functions described in Section 9.26.3, Section 9.26.4, and Section 9.26.5 are also relevant for replication.

#### Table 9-82. Replication SQL Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_create_physical_replication_slot(slot_name, xlog_position pg_lsn)</td>
<td>(slot_name name, xlog_position</td>
<td>Creates a new physical replication slot named slot_name. The optional second parameter, when true, specifies that the LSN for this replication slot be reserved immediately; otherwise the LSN is reserved on first connection from a streaming replication client. Streaming changes from a physical slot is only possible with the streaming-replication protocol — see Section 51.3. This function corresponds to the replication protocol command CREATE_REPLICATION_SLOT ... PHYSICAL.</td>
</tr>
<tr>
<td>pg_drop_replication_slot(slot_name name)</td>
<td>void</td>
<td>Drops the physical or logical replication slot named slot_name. Same as replication protocol command DROP_REPLICATION_SLOT.</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------------------------</td>
<td>------------------------------------------</td>
<td>----------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>pg_create_logical_replication_slot(name, plugin name)</code></td>
<td><code>(slot_name name, xlog_position pg_lsn)</code></td>
<td>Creates a new logical (decoding) replication slot named <code>slot_name</code> using the output plugin <code>plugin</code>. A call to this function has the same effect as the replication protocol command <code>CREATE_REPLICATION_SLOT ... LOGICAL</code>.</td>
</tr>
<tr>
<td><code>pg_logical_slot_get_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</code></td>
<td><code>(location pg_lsn, xid xid, data text)</code></td>
<td>Returns changes in the slot <code>slot_name</code>, starting from the point at which since changes have been consumed last. If <code>upto_lsn</code> and <code>upto_nchanges</code> are NULL, logical decoding will continue until end of WAL. If <code>upto_lsn</code> is non-NULL, decoding will include only those transactions which commit prior to the specified LSN. If <code>upto_nchanges</code> is non-NULL, decoding will stop when the number of rows produced by decoding exceeds the specified value. Note, however, that the actual number of rows returned may be larger, since this limit is only checked after adding the rows produced when decoding each new transaction commit.</td>
</tr>
<tr>
<td><code>pg_logical_slot_peek_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</code></td>
<td><code>(location text, xid xid, data text)</code></td>
<td>Behaves just like the <code>pg_logical_slot_get_changes()</code> function, except that changes are not consumed; that is, they will be returned again on future calls.</td>
</tr>
<tr>
<td><code>pg_logical_slot_get_binary_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</code></td>
<td><code>(location pg_lsn, xid xid, data bytea)</code></td>
<td>Behaves just like the <code>pg_logical_slot_get_changes()</code> function, except that changes are returned as <code>bytea</code>.</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------</td>
<td>-------------</td>
<td>-------------</td>
</tr>
<tr>
<td>pg_logical_slot_peek_binary_changes(slot_name name, upto_lsn pg_lsn, upto_nchanges int, VARIADIC options text[])</td>
<td>(location pg_lsn, xid xid, changes bytea)</td>
<td>Behaves just like the pg_logical_slot_get_changes() function, except that changes are returned as bytea and that changes are not consumed; that is, they will be returned again on future calls.</td>
</tr>
<tr>
<td>pg_replication_origin_create(node_name text)</td>
<td>oid</td>
<td>Create a replication origin with the given external name, and return the internal id assigned to it.</td>
</tr>
<tr>
<td>pg_replication_origin_drop(node_name text)</td>
<td>void</td>
<td>Delete a previously created replication origin, including any associated replay progress.</td>
</tr>
<tr>
<td>pg_replication_origin_oid(node_name text)</td>
<td>oid</td>
<td>Lookup a replication origin by name and return the internal id. If no corresponding replication origin is found an error is thrown.</td>
</tr>
<tr>
<td>pg_replication_origin_session_setup(node_name text)</td>
<td>void</td>
<td>Mark the current session as replaying from the given origin, allowing replay progress to be tracked. Use pg_replication_origin_session_reset to revert. Can only be used if no previous origin is configured.</td>
</tr>
<tr>
<td>pg_replication_origin_session_reset()</td>
<td>void</td>
<td>Cancel the effects of pg_replication_origin_session_setup().</td>
</tr>
<tr>
<td>pg_replication_origin_session_progress(flush bool)</td>
<td>pg_lsn</td>
<td>Return the replay position for the replication origin configured in the current session. The parameter flush determines whether the corresponding local transaction will be guaranteed to have been flushed to disk or not.</td>
</tr>
<tr>
<td>pg_replication_origin_session_is_setup()</td>
<td>bool</td>
<td>Has a replication origin been configured in the current session?</td>
</tr>
<tr>
<td>Function</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------------------------</td>
<td>-------------</td>
<td>------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>pg_replication_origin_xact_setup(origin_lsn pg_lsn, origin_timestamp timestamptz)</code></td>
<td><code>void</code></td>
<td>Mark the current transaction as replaying a transaction that has committed at the given LSN and timestamp. Can only be called when a replication origin has previously been configured using <code>pg_replication_origin_session_setup()</code>.</td>
</tr>
<tr>
<td><code>pg_replication_origin_xact_reset()</code></td>
<td><code>void</code></td>
<td>Cancel the effects of <code>pg_replication_origin_xact_setup()</code>.</td>
</tr>
<tr>
<td><code>pg_replication_origin_advance(node_name text, pos pg_lsn)</code></td>
<td><code>void</code></td>
<td>Set replication progress for the given node to the given position. This primarily is useful for setting up the initial position or a new position after configuration changes and similar. Be aware that careless use of this function can lead to inconsistently replicated data.</td>
</tr>
<tr>
<td><code>pg_replication_origin_progress(node_name text, flush bool)</code></td>
<td><code>pg_lsn</code></td>
<td>Return the replay position for the given replication origin. The parameter <code>flush</code> determines whether the corresponding local transaction will be guaranteed to have been flushed to disk or not.</td>
</tr>
<tr>
<td><code>pg_logical_emit_message(transactional bool, prefix text, content text)</code></td>
<td><code>pg_lsn</code></td>
<td>Emit text logical decoding message. This can be used to pass generic messages to logical decoding plugins through WAL. The parameter <code>transactional</code> specifies if the message should be part of current transaction or if it should be written immediately and decoded as soon as the logical decoding reads the record. The <code>prefix</code> is textual prefix used by the logical decoding plugins to easily recognize interesting messages for them. The <code>content</code> is the text of the message.</td>
</tr>
</tbody>
</table>
Function | Return Type | Description
--- | --- | ---
pg_logical_emit_message(transactional bool, prefix text, content bytes) | pg_lsn | Emit binary logical decoding message. This can be used to pass generic messages to logical decoding plugins through WAL. The parameter transactional specifies if the message should be part of current transaction or if it should be written immediately and decoded as soon as the logical decoding reads the record. The prefix is textual prefix used by the logical decoding plugins to easily recognize interesting messages for them. The content is the binary content of the message.

9.26.7. Database Object Management Functions

The functions shown in Table 9-83 calculate the disk space usage of database objects.

**Table 9-83. Database Object Size Functions**

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_column_size(any)</td>
<td>int</td>
<td>Number of bytes used to store a particular value (possibly compressed)</td>
</tr>
<tr>
<td>pg_database_size(oid)</td>
<td>bigint</td>
<td>Disk space used by the database with the specified OID</td>
</tr>
<tr>
<td>pg_database_size(name)</td>
<td>bigint</td>
<td>Disk space used by the database with the specified name</td>
</tr>
<tr>
<td>pg_indexes_size(regclass)</td>
<td>bigint</td>
<td>Total disk space used by indexes attached to the specified table</td>
</tr>
<tr>
<td>pg_relation_size(relation regclass, fork text)</td>
<td>bigint</td>
<td>Disk space used by the specified fork (&quot;main&quot;, &quot;fsm&quot;, &quot;vm&quot;, or &quot;init&quot;) of the specified table or index</td>
</tr>
<tr>
<td>pg_relation_size(relation regclass)</td>
<td>bigint</td>
<td>Shorthand for pg_relation_size(..., &quot;main&quot;)</td>
</tr>
<tr>
<td>pg_size_bytes(text)</td>
<td>bigint</td>
<td>Converts a size in human-readable format with size units into bytes</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------------</td>
<td>-------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>pg_size_pretty(bigint)</code></td>
<td>text</td>
<td>Converts a size in bytes expressed as a 64-bit integer into a human-readable format with size units</td>
</tr>
<tr>
<td><code>pg_size_pretty(numeric)</code></td>
<td>text</td>
<td>Converts a size in bytes expressed as a numeric value into a human-readable format with size units</td>
</tr>
<tr>
<td><code>pg_table_size(regclass)</code></td>
<td>bigint</td>
<td>Disk space used by the specified table, excluding indexes (but including TOAST, free space map, and visibility map)</td>
</tr>
<tr>
<td><code>pg_tablespace_size(oid)</code></td>
<td>bigint</td>
<td>Disk space used by the tablespace with the specified OID</td>
</tr>
<tr>
<td><code>pg_tablespace_size(name)</code></td>
<td>bigint</td>
<td>Disk space used by the tablespace with the specified name</td>
</tr>
<tr>
<td><code>pg_total_relation_size(regclass)</code></td>
<td>bigint</td>
<td>Total disk space used by the specified table, including all indexes and TOAST data</td>
</tr>
</tbody>
</table>

`pg_column_size` shows the space used to store any individual data value.

`pg_total_relation_size` accepts the OID or name of a table or toast table, and returns the total on-disk space used for that table, including all associated indexes. This function is equivalent to `pg_table_size + pg_indexes_size`.

`pg_table_size` accepts the OID or name of a table and returns the disk space needed for that table, exclusive of indexes. (TOAST space, free space map, and visibility map are included.)

`pg_indexes_size` accepts the OID or name of a table and returns the total disk space used by all the indexes attached to that table.

`pg_database_size` and `pg_tablespace_size` accept the OID or name of a database or tablespace, and return the total disk space used therein. To use `pg_database_size`, you must have CONNECT permission on the specified database (which is granted by default). To use `pg_tablespace_size`, you must have CREATE permission on the specified tablespace, unless it is the default tablespace for the current database.

`pg_relation_size` accepts the OID or name of a table, index or toast table, and returns the on-disk size in bytes of one fork of that relation. (Note that for most purposes it is more convenient to use the higher-level functions `pg_total_relation_size` or `pg_table_size`, which sum the sizes of all forks.) With one argument, it returns the size of the main data fork of the relation. The second argument can be provided to specify which fork to examine:

- ‘main’ returns the size of the main data fork of the relation.
- ‘fsm’ returns the size of the Free Space Map (see Section 65.3) associated with the relation.
- ‘vm’ returns the size of the Visibility Map (see Section 65.4) associated with the relation.
- ‘init’ returns the size of the initialization fork, if any, associated with the relation.
Chapter 9. Functions and Operators

`pg_size_pretty` can be used to format the result of one of the other functions in a human-readable way, using bytes, kB, MB, GB or TB as appropriate.

`pg_size_bytes` can be used to get the size in bytes from a string in human-readable format. The input may have units of bytes, kB, MB, GB or TB, and is parsed case-insensitively. If no units are specified, bytes are assumed.

**Note:** The units kB, MB, GB and TB used by the functions `pg_size_pretty` and `pg_size_bytes` are defined using powers of 2 rather than powers of 10, so 1kB is 1024 bytes, 1MB is 1024² = 1048576 bytes, and so on.

The functions above that operate on tables or indexes accept a `regclass` argument, which is simply the OID of the table or index in the `pg_class` system catalog. You do not have to look up the OID by hand, however, since the `regclass` data type’s input converter will do the work for you. Just write the table name enclosed in single quotes so that it looks like a literal constant. For compatibility with the handling of ordinary SQL names, the string will be converted to lower case unless it contains double quotes around the table name.

If an OID that does not represent an existing object is passed as argument to one of the above functions, NULL is returned.

The functions shown in Table 9-84 assist in identifying the specific disk files associated with database objects.

**Table 9-84. Database Object Location Functions**

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_relation_filenode</code></td>
<td>oid</td>
<td>Filenode number of the specified relation</td>
</tr>
<tr>
<td><code>pg_relation_filepath</code></td>
<td>text</td>
<td>File path name of the specified relation</td>
</tr>
<tr>
<td><code>pg_filenode_relation</code></td>
<td>regclass</td>
<td>Find the relation associated with a given tablespace and filenode</td>
</tr>
</tbody>
</table>

`pg_relation_filenode` accepts the OID or name of a table, index, sequence, or toast table, and returns the “filenode” number currently assigned to it. The filenode is the base component of the file name(s) used for the relation (see Section 65.1 for more information). For most tables the result is the same as `pg_class.relfilenode`, but for certain system catalogs `relfilenode` is zero and this function must be used to get the correct value. The function returns NULL if passed a relation that does not have storage, such as a view.

`pg_relation_filepath` is similar to `pg_relation_filenode`, but it returns the entire file path name (relative to the database cluster’s data directory `PGDATA`) of the relation.

`pg_filenode_relation` is the reverse of `pg_relation_filenode`. Given a “tablespace” OID and a “filenode”, it returns the associated relation’s OID. For a table in the database’s default tablespace, the tablespace can be specified as 0.
9.26.8. Index Maintenance Functions

Table 9-85 shows the functions available for index maintenance tasks. These functions cannot be executed during recovery. Use of these functions is restricted to superusers and the owner of the given index.

### Table 9-85. Index Maintenance Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>brin_summarize_new_values(index regclass)</code></td>
<td>integer</td>
<td>summarize page ranges not already summarized</td>
</tr>
<tr>
<td><code>gin_clean_pending_list(index regclass)</code></td>
<td>bigint</td>
<td>move GIN pending list entries into main index structure</td>
</tr>
</tbody>
</table>

`brin_summarize_new_values` accepts the OID or name of a BRIN index and inspects the index to find page ranges in the base table that are not currently summarized by the index; for any such range it creates a new summary index tuple by scanning the table pages. It returns the number of new page range summaries that were inserted into the index.

`gin_clean_pending_list` accepts the OID or name of a GIN index and cleans up the pending list of the specified index by moving entries in it to the main GIN data structure in bulk. It returns the number of pages removed from the pending list. Note that if the argument is a GIN index built with the `fastupdate` option disabled, no cleanup happens and the return value is 0, because the index doesn’t have a pending list. Please see Section 63.4.1 and Section 63.5 for details of the pending list and `fastupdate` option.


The functions shown in Table 9-86 provide native access to files on the machine hosting the server. Only files within the database cluster directory and the `log_directory` can be accessed. Use a relative path for files in the cluster directory, and a path matching the `log_directory` configuration setting for log files. Use of these functions is restricted to superusers.

### Table 9-86. Generic File Access Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_ls_dir(dirname text [, missing_ok boolean, include_dot_dirs boolean])</code></td>
<td>setof text</td>
<td>List the contents of a directory.</td>
</tr>
<tr>
<td><code>pg_read_file(filename text [, offset bigint, length bigint [, missing_ok boolean ] ])</code></td>
<td>text</td>
<td>Return the contents of a text file.</td>
</tr>
<tr>
<td><code>pg_read_binary_file(filename text [, offset bigint, length bigint [, missing_ok boolean ] ])</code></td>
<td>bytea</td>
<td>Return the contents of a file.</td>
</tr>
</tbody>
</table>
9.26.10. Advisory Lock Functions

The functions shown in Table 9-87 manage advisory locks. For details about proper use of these functions, see Section 13.3.5.

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_advisory_lock(key bigint)</td>
<td>void</td>
<td>Obtain exclusive session level advisory lock</td>
</tr>
<tr>
<td>pg_advisory_lock(key1 int, key2 int)</td>
<td>void</td>
<td>Obtain exclusive session level advisory lock</td>
</tr>
<tr>
<td>pg_advisory_lock_shared(key bigint)</td>
<td>void</td>
<td>Obtain shared session level advisory lock</td>
</tr>
<tr>
<td>pg_advisory_lock_shared(key1 int, key2 int)</td>
<td>void</td>
<td>Obtain shared session level advisory lock</td>
</tr>
<tr>
<td>Name</td>
<td>Return Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------------------------------</td>
<td>-------------</td>
<td>-------------</td>
</tr>
<tr>
<td><code>pg_advisory_unlock(key bigint)</code></td>
<td>boolean</td>
<td>Release an exclusive session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_unlock(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Release an exclusive session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_unlock_all()</code></td>
<td>void</td>
<td>Release all session level advisory locks held by the current session</td>
</tr>
<tr>
<td><code>pg_advisory_unlock_shared(key bigint)</code></td>
<td>boolean</td>
<td>Release a shared session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_unlock_shared(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Release a shared session level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_xact_lock(key bigint)</code></td>
<td>void</td>
<td>Obtain exclusive transaction level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_xact_lock(key1 int, key2 int)</code></td>
<td>void</td>
<td>Obtain exclusive transaction level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_xact_lock_shared(key bigint)</code></td>
<td>void</td>
<td>Obtain shared transaction level advisory lock</td>
</tr>
<tr>
<td><code>pg_advisory_xact_lock_shared(key1 int, key2 int)</code></td>
<td>void</td>
<td>Obtain shared transaction level advisory lock</td>
</tr>
<tr>
<td><code>pg_try_advisory_lock(key bigint)</code></td>
<td>boolean</td>
<td>Obtain exclusive session level advisory lock if available</td>
</tr>
<tr>
<td><code>pg_try_advisory_lock(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Obtain exclusive session level advisory lock if available</td>
</tr>
<tr>
<td><code>pg_try_advisory_lock_shared(key bigint)</code></td>
<td>boolean</td>
<td>Obtain shared session level advisory lock if available</td>
</tr>
<tr>
<td><code>pg_try_advisory_lock_shared(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Obtain shared session level advisory lock if available</td>
</tr>
<tr>
<td><code>pg_try_advisory_xact_lock(key bigint)</code></td>
<td>boolean</td>
<td>Obtain exclusive transaction level advisory lock if available</td>
</tr>
<tr>
<td><code>pg_try_advisory_xact_lock(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Obtain exclusive transaction level advisory lock if available</td>
</tr>
<tr>
<td><code>pg_try_advisory_xact_lock_shared(key bigint)</code></td>
<td>boolean</td>
<td>Obtain shared transaction level advisory lock if available</td>
</tr>
<tr>
<td><code>pg_try_advisory_xact_lock_shared(key1 int, key2 int)</code></td>
<td>boolean</td>
<td>Obtain shared transaction level advisory lock if available</td>
</tr>
</tbody>
</table>

`pg_advisory_lock` locks an application-defined resource, which can be identified either by a single
64-bit key value or two 32-bit key values (note that these two key spaces do not overlap). If another session already holds a lock on the same resource identifier, this function will wait until the resource becomes available. The lock is exclusive. Multiple lock requests stack, so that if the same resource is locked three times it must then be unlocked three times to be released for other sessions’ use.

`pg_advisory_lock_shared` works the same as `pg_advisory_lock`, except the lock can be shared with other sessions requesting shared locks. Only would-be exclusive lockers are locked out.

`pg_try_advisory_lock` is similar to `pg_advisory_lock`, except the function will not wait for the lock to become available. It will either obtain the lock immediately and return `true`, or return `false` if the lock cannot be acquired immediately.

`pg_try_advisory_lock_shared` works the same as `pg_try_advisory_lock`, except it attempts to acquire a shared rather than an exclusive lock.

`pg_advisory_unlock` will release a previously-acquired exclusive session level advisory lock. It returns `true` if the lock is successfully released. If the lock was not held, it will return `false`, and in addition, an SQL warning will be reported by the server.

`pg_advisory_unlock_shared` works the same as `pg_advisory_unlock`, except it releases a shared session level advisory lock.

`pg_advisory_unlock_all` will release all session level advisory locks held by the current session. (This function is implicitly invoked at session end, even if the client disconnects ungracefully.)

`pg_advisory_xact_lock` works the same as `pg_advisory_lock`, except the lock is automatically released at the end of the current transaction and cannot be released explicitly.

`pg_advisory_xact_lock_shared` works the same as `pg_advisory_lock_shared`, except the lock is automatically released at the end of the current transaction and cannot be released explicitly.

`pg_try_advisory_xact_lock` works the same as `pg_try_advisory_lock`, except the lock, if acquired, is automatically released at the end of the current transaction and cannot be released explicitly.

`pg_try_advisory_xact_lock_shared` works the same as `pg_try_advisory_lock_shared`, except the lock, if acquired, is automatically released at the end of the current transaction and cannot be released explicitly.

### 9.27. Trigger Functions

Currently PostgreSQL provides one built-in trigger function, `suppress_redundant_updates_trigger`, which will prevent any update that does not actually change the data in the row from taking place, in contrast to the normal behavior which always performs the update regardless of whether or not the data has changed. (This normal behavior makes updates run faster, since no checking is required, and is also useful in certain cases.) Ideally, you should normally avoid running updates that don’t actually change the data in the record. Redundant updates can cost considerable unnecessary time, especially if there are lots of indexes to alter, and space in dead rows that will eventually have to be vacuumed. However, detecting such situations in client code is not always easy, or even possible, and writing expressions to detect them can be error-prone. An alternative is to use `suppress_redundant_updates_trigger`, which will skip updates that don’t change the data. You should use this with care, however. The trigger takes a small but non-trivial time for each record, so if most of the records affected by an update are actually changed, use of this trigger will actually make the update run slower.

The `suppress_redundant_updates_trigger` function can be added to a table like this:
CREATE TRIGGER z_min_update  
BEFORE UPDATE ON tablename  
FOR EACH ROW EXECUTE PROCEDURE suppress_redundant_updates_trigger();

In most cases, you would want to fire this trigger last for each row. Bearing in mind that triggers fire in name order, you would then choose a trigger name that comes after the name of any other trigger you might have on the table.

For more information about creating triggers, see CREATE TRIGGER.

9.28. Event Trigger Functions

PostgreSQL provides these helper functions to retrieve information from event triggers.

For more information about event triggers, see Chapter 38.

9.28.1. Capturing Changes at Command End

`pg_event_trigger_ddl_commands` returns a list of DDL commands executed by each user action, when invoked in a function attached to a `ddl_command_end` event trigger. If called in any other context, an error is raised. `pg_event_trigger_ddl_commands` returns one row for each base command executed; some commands that are a single SQL sentence may return more than one row. This function returns the following columns:

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>classid</td>
<td>oid</td>
<td>OID of catalog the object belongs in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>OID of the object itself</td>
</tr>
<tr>
<td>objsubid</td>
<td>integer</td>
<td>Sub-object ID (e.g. attribute number for a column)</td>
</tr>
<tr>
<td>command_tag</td>
<td>text</td>
<td>Command tag</td>
</tr>
<tr>
<td>object_type</td>
<td>text</td>
<td>Type of the object</td>
</tr>
<tr>
<td>schema_name</td>
<td>text</td>
<td>Name of the schema the object belongs in, if any; otherwise NULL. No quoting is applied.</td>
</tr>
<tr>
<td>object_identity</td>
<td>text</td>
<td>Text rendering of the object identity, schema-qualified. Each identifier included in the identity is quoted if necessary.</td>
</tr>
<tr>
<td>in_extension</td>
<td>bool</td>
<td>True if the command is part of an extension script</td>
</tr>
<tr>
<td>command</td>
<td><code>pg_ddl_command</code></td>
<td>A complete representation of the command, in internal format. This cannot be output directly, but it can be passed to other functions to obtain different pieces of information about the command.</td>
</tr>
</tbody>
</table>
9.28.2. Processing Objects Dropped by a DDL Command

`pg_event_trigger_dropped_objects` returns a list of all objects dropped by the command in whose `sql_drop` event it is called. If called in any other context, `pg_event_trigger_dropped_objects` raises an error.

`pg_event_trigger_dropped_objects` returns the following columns:

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>classid</td>
<td>oid</td>
<td>OID of catalog the object belonged in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>OID of the object itself</td>
</tr>
<tr>
<td>objsubid</td>
<td>integer</td>
<td>Sub-object ID (e.g. attribute number for a column)</td>
</tr>
<tr>
<td>original</td>
<td>bool</td>
<td>True if this was one of the root object(s) of the deletion</td>
</tr>
<tr>
<td>normal</td>
<td>bool</td>
<td>True if there was a normal dependency relationship in the dependency graph leading to this object</td>
</tr>
<tr>
<td>is_temporary</td>
<td>bool</td>
<td>True if this was a temporary object</td>
</tr>
<tr>
<td>object_type</td>
<td>text</td>
<td>Type of the object</td>
</tr>
<tr>
<td>schema_name</td>
<td>text</td>
<td>Name of the schema the object belonged in, if any; otherwise NULL. No quoting is applied.</td>
</tr>
<tr>
<td>object_name</td>
<td>text</td>
<td>Name of the object, if the combination of schema and name can be used as a unique identifier for the object; otherwise NULL. No quoting is applied, and name is never schema-qualified.</td>
</tr>
<tr>
<td>object_identity</td>
<td>text</td>
<td>Text rendering of the object identity, schema-qualified. Each identifier included in the identity is quoted if necessary.</td>
</tr>
<tr>
<td>address_names</td>
<td>text[]</td>
<td>An array that, together with <code>object_type</code> and <code>address_args</code>, can be used by the <code>pg_get_object_address()</code> function to recreate the object address in a remote server containing an identically named object of the same kind</td>
</tr>
</tbody>
</table>
The \texttt{pg\_event\_trigger\_dropped\_objects} function can be used in an event trigger like this:

```sql
CREATE FUNCTION test_event_trigger_for_drops()
RETURNS event_trigger LANGUAGE plpgsql AS $$
DECLARE
obj record;
BEGIN
FOR obj IN SELECT * FROM pg_event_trigger_dropped_objects()
 LOOP
RAISE NOTICE '% dropped object: % % % %',
tg_tag,
obj.object_type,
obj.schema_name,
obj.object_name,
obj.object_identity;
END LOOP;
END LOOP;
END
$$;
CREATE EVENT TRIGGER test_event_trigger_for_drops
ON sql_drop
EXECUTE PROCEDURE test_event_trigger_for_drops();
```

### 9.28.3. Handling a Table Rewrite Event

The functions shown in Table 9-88 provide information about a table for which a \texttt{table\_rewrite} event has just been called. If called in any other context, an error is raised.

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{pg_event_trigger_table_rewrite_oid()}</td>
<td>Oid</td>
<td>The OID of the table about to be rewritten.</td>
</tr>
<tr>
<td>\texttt{pg_event_trigger_table_rewrite_reason()}</td>
<td>int</td>
<td>The reason code(s) explaining the reason for rewriting. The exact meaning of the codes is release dependent.</td>
</tr>
</tbody>
</table>

The \texttt{pg\_event\_trigger\_table\_rewrite\_oid} function can be used in an event trigger like this:

```sql
CREATE FUNCTION test_event_trigger_table_rewrite_oid()
RETURNS event_trigger LANGUAGE plpgsql AS $$
BEGIN
RAISE NOTICE 'rewriting table % for reason %',
```
pg_event_trigger_table_rewrite_oid(): regclass,
pg_event_trigger_table_rewrite_reason();

END;
$$;

CREATE EVENT TRIGGER test_table_rewrite_oid
  ON table_rewrite
  EXECUTE PROCEDURE test_event_trigger_table_rewrite_oid();
Chapter 10. Type Conversion

SQL statements can, intentionally or not, require the mixing of different data types in the same expression. PostgreSQL has extensive facilities for evaluating mixed-type expressions.

In many cases a user does not need to understand the details of the type conversion mechanism. However, implicit conversions done by PostgreSQL can affect the results of a query. When necessary, these results can be tailored by using explicit type conversion.

This chapter introduces the PostgreSQL type conversion mechanisms and conventions. Refer to the relevant sections in Chapter 8 and Chapter 9 for more information on specific data types and allowed functions and operators.

10.1. Overview

SQL is a strongly typed language. That is, every data item has an associated data type which determines its behavior and allowed usage. PostgreSQL has an extensible type system that is more general and flexible than other SQL implementations. Hence, most type conversion behavior in PostgreSQL is governed by general rules rather than by ad hoc heuristics. This allows the use of mixed-type expressions even with user-defined types.

The PostgreSQL scanner/parser divides lexical elements into five fundamental categories: integers, non-integer numbers, strings, identifiers, and key words. Constants of most non-numeric types are first classified as strings. The SQL language definition allows specifying type names with strings, and this mechanism can be used in PostgreSQL to start the parser down the correct path. For example, the query:

```
SELECT text 'Origin' AS "label", point '(0,0)' AS "value";
```

| label | value |
|-------+-------|
| Origin | (0,0) |

has two literal constants, of type `text` and `point`. If a type is not specified for a string literal, then the placeholder type `unknown` is assigned initially, to be resolved in later stages as described below.

There are four fundamental SQL constructs requiring distinct type conversion rules in the PostgreSQL parser:

Function calls

Much of the PostgreSQL type system is built around a rich set of functions. Functions can have one or more arguments. Since PostgreSQL permits function overloading, the function name alone does not uniquely identify the function to be called; the parser must select the right function based on the data types of the supplied arguments.

Operators

PostgreSQL allows expressions with prefix and postfix unary (one-argument) operators, as well as binary (two-argument) operators. Like functions, operators can be overloaded, so the same problem of selecting the right operator exists.
Chapter 10. Type Conversion

Value Storage

SQL `INSERT` and `UPDATE` statements place the results of expressions into a table. The expressions in the statement must be matched up with, and perhaps converted to, the types of the target columns.

`UNION`, `CASE`, and related constructs

Since all query results from a unionized `SELECT` statement must appear in a single set of columns, the types of the results of each `SELECT` clause must be matched up and converted to a uniform set. Similarly, the result expressions of a `CASE` construct must be converted to a common type so that the `CASE` expression as a whole has a known output type. The same holds for `ARRAY` constructs, and for the `GREATEST` and `LEAST` functions.

The system catalogs store information about which conversions, or castings, exist between which data types, and how to perform those conversions. Additional casts can be added by the user with the `CREATE CAST` command. (This is usually done in conjunction with defining new data types. The set of casts between built-in types has been carefully crafted and is best not altered.)

An additional heuristic provided by the parser allows improved determination of the proper casting behavior among groups of types that have implicit casts. Data types are divided into several basic type categories, including boolean, numeric, string, bitstring, datetime, timespan, geometric, network, and user-defined. (For a list see Table 50-56; but note it is also possible to create custom type categories.) Within each category there can be one or more preferred types, which are preferred when there is a choice of possible types. With careful selection of preferred types and available implicit casts, it is possible to ensure that ambiguous expressions (those with multiple candidate parsing solutions) can be resolved in a useful way.

All type conversion rules are designed with several principles in mind:

- Implicit conversions should never have surprising or unpredictable outcomes.
- There should be no extra overhead in the parser or executor if a query does not need implicit type conversion. That is, if a query is well-formed and the types already match, then the query should execute without spending extra time in the parser and without introducing unnecessary implicit conversion calls in the query.
- Additionally, if a query usually requires an implicit conversion for a function, and if then the user defines a new function with the correct argument types, the parser should use this new function and no longer do implicit conversion to use the old function.

10.2. Operators

The specific operator that is referenced by an operator expression is determined using the following procedure. Note that this procedure is indirectly affected by the precedence of the operators involved, since that will determine which sub-expressions are taken to be the inputs of which operators. See Section 4.1.6 for more information.

Operator Type Resolution

1. Select the operators to be considered from the `pg_operator` system catalog. If a non-schema-qualified operator name was used (the usual case), the operators considered are those with the
Chapter 10. Type Conversion

matching name and argument count that are visible in the current search path (see Section 5.8.3). If a qualified operator name was given, only operators in the specified schema are considered.

a. If the search path finds multiple operators with identical argument types, only the one appearing earliest in the path is considered. Operators with different argument types are considered on an equal footing regardless of search path position.

2. Check for an operator accepting exactly the input argument types. If one exists (there can be only one exact match in the set of operators considered), use it. Lack of an exact match creates a security hazard when calling, via qualified name ¹ (not typical), any operator found in a schema that permits untrusted users to create objects. In such situations, cast arguments to force an exact match.

a. If one argument of a binary operator invocation is of the unknown type, then assume it is the same type as the other argument for this check. Invocations involving two unknown inputs, or a unary operator with an unknown input, will never find a match at this step.

b. If one argument of a binary operator invocation is of the unknown type and the other is of a domain type, next check to see if there is an operator accepting exactly the domain’s base type on both sides; if so, use it.

3. Look for the best match.

a. Discard candidate operators for which the input types do not match and cannot be converted (using an implicit conversion) to match. unknown literals are assumed to be convertible to anything for this purpose. If only one candidate remains, use it; else continue to the next step.

b. If any input argument is of a domain type, treat it as being of the domain’s base type for all subsequent steps. This ensures that domains act like their base types for purposes of ambiguous-operator resolution.

c. Run through all candidates and keep those with the most exact matches on input types. Keep all candidates if none have exact matches. If only one candidate remains, use it; else continue to the next step.

d. Run through all candidates and keep those that accept preferred types (of the input data type’s type category) at the most positions where type conversion will be required. Keep all candidates if none accept preferred types. If only one candidate remains, use it; else continue to the next step.

e. If any input arguments are unknown, check the type categories accepted at those argument positions by the remaining candidates. At each position, select the string category if any candidate accepts that category. (This bias towards string is appropriate since an unknown-type literal looks like a string.) Otherwise, if all the remaining candidates accept the same type category, select that category; otherwise fail because the correct choice cannot be deduced without more clues. Now discard candidates that do not accept the selected type category. Furthermore, if any candidate accepts a preferred type in that category, discard candidates that accept non-preferred types for that argument. Keep all candidates if none survive these tests. If only one candidate remains, use it; else continue to the next step.

f. If there are both unknown and known-type arguments, and all the known-type arguments have the same type, assume that the unknown arguments are also of that type.

¹ The hazard does not arise with a non-schema-qualified name, because a search path containing schemas that permit untrusted users to create objects is not a secure schema usage pattern.
Chapter 10. Type Conversion

and check which candidates can accept that type at the unknown-argument positions. If exactly one candidate passes this test, use it. Otherwise, fail.

Some examples follow.

**Example 10-1. Factorial Operator Type Resolution**

There is only one factorial operator (postfix !) defined in the standard catalog, and it takes an argument of type bigint. The scanner assigns an initial type of integer to the argument in this query expression:

```
SELECT 40 ! AS "40 factorial";
```

```
40 factorial
--------------
8159152832478977343456112695961158942720000000000
(1 row)
```

So the parser does a type conversion on the operand and the query is equivalent to:

```
SELECT CAST(40 AS bigint) ! AS "40 factorial";
```

**Example 10-2. String Concatenation Operator Type Resolution**

A string-like syntax is used for working with string types and for working with complex extension types. Strings with unspecified type are matched with likely operator candidates.

An example with one unspecified argument:

```
SELECT text 'abc' || 'def' AS "text and unknown";
```

```
text and unknown
---------------
abcdef
(1 row)
```

In this case the parser looks to see if there is an operator taking text for both arguments. Since there is, it assumes that the second argument should be interpreted as type text.

Here is a concatenation of two values of unspecified types:

```
SELECT 'abc' || 'def' AS "unspecified";
```

```
unspecified
----------
abcdef
(1 row)
```

In this case there is no initial hint for which type to use, since no types are specified in the query. So, the parser looks for all candidate operators and finds that there are candidates accepting both string-category and bit-string-category inputs. Since string category is preferred when available, that category is selected, and then the preferred type for strings, text, is used as the specific type to resolve the unknown-type literals as.
Example 10-3. Absolute-Value and Negation Operator Type Resolution

The PostgreSQL operator catalog has several entries for the prefix operator @, all of which implement absolute-value operations for various numeric data types. One of these entries is for type float8, which is the preferred type in the numeric category. Therefore, PostgreSQL will use that entry when faced with an unknown input:

```
SELECT @ '-4.5' AS "abs";
```

```
  abs
----
  4.5
```

(1 row)

Here the system has implicitly resolved the unknown-type literal as type float8 before applying the chosen operator. We can verify that float8 and not some other type was used:

```
SELECT @ '-4.5e500' AS "abs";
```

ERROR: "-4.5e500" is out of range for type double precision

On the other hand, the prefix operator ~ (bitwise negation) is defined only for integer data types, not for float8. So, if we try a similar case with ~, we get:

```
SELECT ~ '20' AS "negation";
```

ERROR: operator is not unique: ~ "unknown"

HINT: Could not choose a best candidate operator. You might need to add explicit type casts.

This happens because the system cannot decide which of the several possible ~ operators should be preferred. We can help it out with an explicit cast:

```
SELECT ~ CAST('20' AS int8) AS "negation";
```

```
  negation
----------
   -21
```

(1 row)

Example 10-4. Array Inclusion Operator Type Resolution

Here is another example of resolving an operator with one known and one unknown input:

```
SELECT array[1,2] <@ '{1,2,3}' as "is subset";
```

```
  is subset
----------
    t
```

(1 row)

The PostgreSQL operator catalog has several entries for the infix operator <@, but the only two that could possibly accept an integer array on the left-hand side are array inclusion (anyarray <@ anyarray) and range inclusion (anymelement <@ anyrange). Since none of these polymorphic pseudo-types (see Section 8.20) are considered preferred, the parser cannot resolve the ambiguity on that basis. However, step 3.f tells it to assume that the unknown-type literal is of the same type as the other input, that is, integer array. Now only one of the two operators can match, so array inclusion is selected. (Had range inclusion been selected, we would have gotten an error, because the string does not have the right format to be a range literal.)
Example 10-5. Custom Operator on a Domain Type

Users sometimes try to declare operators applying just to a domain type. This is possible but is not nearly as useful as it might seem, because the operator resolution rules are designed to select operators applying to the domain’s base type. As an example consider

```sql
CREATE DOMAIN mytext AS text CHECK(...);
CREATE FUNCTION mytext_eq_text (mytext, text) RETURNS boolean AS ...;
CREATE OPERATOR = (procedure=mytext_eq_text, leftarg=mytext, rightarg=text);
CREATE TABLE mytable (val mytext);
SELECT * FROM mytable WHERE val = 'foo';
```

This query will not use the custom operator. The parser will first see if there is a `mytext = mytext` operator (step 2.a), which there is not; then it will consider the domain’s base type `text`, and see if there is a `text = text` operator (step 2.b), which there is; so it resolves the unknown-type literal as `text` and uses the `text = text` operator. The only way to get the custom operator to be used is to explicitly cast the literal:

```sql
SELECT * FROM mytable WHERE val = text 'foo';
```

so that the `mytext = text` operator is found immediately according to the exact-match rule. If the best-match rules are reached, they actively discriminate against operators on domain types. If they did not, such an operator would create too many ambiguous-operator failures, because the casting rules always consider a domain as castable to or from its base type, and so the domain operator would be considered usable in all the same cases as a similarly-named operator on the base type.

10.3. Functions

The specific function that is referenced by a function call is determined using the following procedure.

**Function Type Resolution**

1. Select the functions to be considered from the `pg_proc` system catalog. If a non-schema-qualified function name was used, the functions considered are those with the matching name and argument count that are visible in the current search path (see Section 5.8.3). If a qualified function name was given, only functions in the specified schema are considered.

   a. If the search path finds multiple functions of identical argument types, only the one appearing earliest in the path is considered. Functions of different argument types are considered on an equal footing regardless of search path position.

   b. If a function is declared with a `VARIADIC` array parameter, and the call does not use the `VARIADIC` keyword, then the function is treated as if the array parameter were replaced by one or more occurrences of its element type, as needed to match the call. After such expansion the function might have effective argument types identical to some non-variadic function. In that case the function appearing earlier in the search path is used, or if the two functions are in the same schema, the non-variadic one is preferred.

   This creates a security hazard when calling, via qualified name \(^2\), a variadic function found in a schema that permits untrusted users to create objects. A malicious user can

---

\(^2\) The hazard does not arise with a non-schema-qualified name, because a search path containing schemas that permit untrusted users to create objects is not a secure schema usage pattern.
take control and execute arbitrary SQL functions as though you executed them. Substitute a call bearing the **VARIADIC** keyword, which bypasses this hazard. Calls populating **VARIADIC "any"** parameters often have no equivalent formulation containing the **VARIADIC** keyword. To issue those calls safely, the function’s schema must permit only trusted users to create objects.

c. Functions that have default values for parameters are considered to match any call that omits zero or more of the defaultable parameter positions. If more than one such function matches a call, the one appearing earliest in the search path is used. If there are two or more such functions in the same schema with identical parameter types in the non-defaulted positions (which is possible if they have different sets of defaultable parameters), the system will not be able to determine which to prefer, and so an “ambiguous function call” error will result if no better match to the call can be found.

This creates an availability hazard when calling, via qualified name\(^2\), any function found in a schema that permits untrusted users to create objects. A malicious user can create a function with the name of an existing function, replicating that function’s parameters and appending novel parameters having default values. This precludes new calls to the original function. To forestall this hazard, place functions in schemas that permit only trusted users to create objects.

2. Check for a function accepting exactly the input argument types. If one exists (there can be only one exact match in the set of functions considered), use it. Lack of an exact match creates a security hazard when calling, via qualified name\(^2\), a function found in a schema that permits untrusted users to create objects. In such situations, cast arguments to force an exact match. (Cases involving **unknown** will never find a match at this step.)

3. If no exact match is found, see if the function call appears to be a special type conversion request. This happens if the function call has just one argument and the function name is the same as the (internal) name of some data type. Furthermore, the function argument must be either an unknown-type literal, or a type that is binary-coercible to the named data type, or a type that could be converted to the named data type by applying that type’s I/O functions (that is, the conversion is either to or from one of the standard string types). When these conditions are met, the function call is treated as a form of **CAST** specification.\(^3\)

4. Look for the best match.

   a. Discard candidate functions for which the input types do not match and cannot be converted (using an implicit conversion) to match. **unknown** literals are assumed to be convertible to anything for this purpose. If only one candidate remains, use it; else continue to the next step.

   b. If any input argument is of a domain type, treat it as being of the domain’s base type for all subsequent steps. This ensures that domains act like their base types for purposes of ambiguous-function resolution.

   c. Run through all candidates and keep those with the most exact matches on input types. Keep all candidates if none have exact matches. If only one candidate remains, use it; else continue to the next step.

---

\(^2\) The hazard does not arise with a non-schema-qualified name, because a search path containing schemas that permit untrusted users to create objects is not a secure schema usage pattern.

\(^2\) The hazard does not arise with a non-schema-qualified name, because a search path containing schemas that permit untrusted users to create objects is not a secure schema usage pattern.

\(^3\) The reason for this step is to support function-style cast specifications in cases where there is not an actual cast function. If there is a cast function, it is conventionally named after its output type, and so there is no need to have a special case. See CREATE CAST for additional commentary.

354
d. Run through all candidates and keep those that accept preferred types (of the input data type’s type category) at the most positions where type conversion will be required. Keep all candidates if none accept preferred types. If only one candidate remains, use it; else continue to the next step.

e. If any input arguments are unknown, check the type categories accepted at those argument positions by the remaining candidates. At each position, select the string category if any candidate accepts that category. (This bias towards string is appropriate since an unknown-type literal looks like a string.) Otherwise, if all the remaining candidates accept the same type category, select that category; otherwise fail because the correct choice cannot be deduced without more clues. Now discard candidates that do not accept the selected type category. Furthermore, if any candidate accepts a preferred type in that category, discard candidates that accept non-preferred types for that argument. Keep all candidates if none survive these tests. If only one candidate remains, use it; else continue to the next step.

f. If there are both unknown and known-type arguments, and all the known-type arguments have the same type, assume that the unknown arguments are also of that type, and check which candidates can accept that type at the unknown-argument positions. If exactly one candidate passes this test, use it. Otherwise, fail.

Note that the “best match” rules are identical for operator and function type resolution. Some examples follow.

Example 10-6. Rounding Function Argument Type Resolution

There is only one round function that takes two arguments; it takes a first argument of type numeric and a second argument of type integer. So the following query automatically converts the first argument of type integer to numeric:

```
SELECT round(4, 4);
```

```
round
--------
4.0000
(1 row)
```

That query is actually transformed by the parser to:

```
SELECT round(CAST (4 AS numeric), 4);
```

Since numeric constants with decimal points are initially assigned the type numeric, the following query will require no type conversion and therefore might be slightly more efficient:

```
SELECT round(4.0, 4);
```

Example 10-7. Variadic Function Resolution

```
CREATE FUNCTION public.variadic_example(VARIADIC numeric[]) RETURNS int
    LANGUAGE sql AS 'SELECT 1';
```

This function accepts, but does not require, the VARIADIC keyword. It tolerates both integer and numeric arguments:

```
SELECT public.variadic_example(0),
    public.variadic_example(0.0),
    public.variadic_example(VARIADIC array[0.0]);
```
However, the first and second calls will prefer more-specific functions, if available:

```sql
CREATE FUNCTION public.variadic_example(numeric) RETURNS int
    LANGUAGE sql AS 'SELECT 2';
CREATE FUNCTION public.variadic_example(int) RETURNS int
    LANGUAGE sql AS 'SELECT 3';
```

```
SELECT public.variadic_example(0),
    public.variadic_example(0.0),
    public.variadic_example(VARIADIC array[0.0]);
```

<table>
<thead>
<tr>
<th>variadic_example</th>
<th>variadic_example</th>
<th>variadic_example</th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>2</td>
<td>1</td>
</tr>
</tbody>
</table>

(1 row)

Given the default configuration and only the first function existing, the first and second calls are insecure. Any user could intercept them by creating the second or third function. By matching the argument type exactly and using the `VARIADIC` keyword, the third call is secure.

---

### Example 10-8. Substring Function Type Resolution

There are several `substr` functions, one of which takes types `text` and `integer`. If called with a string constant of unspecified type, the system chooses the candidate function that accepts an argument of the preferred category `string` (namely of type `text`).

```sql
SELECT substr('1234', 3);
```

```
substr
-------
  34
(1 row)
```

If the string is declared to be of type `varchar`, as might be the case if it comes from a table, then the parser will try to convert it to become `text`:

```sql
SELECT substr(varchar '1234', 3);
```

```
substr
-------
  34
(1 row)
```

This is transformed by the parser to effectively become:

```sql
SELECT substr(CAST (varchar '1234' AS text), 3);
```

**Note:** The parser learns from the `pg_cast` catalog that `text` and `varchar` are binary-compatible, meaning that one can be passed to a function that accepts the other without doing any physical conversion. Therefore, no type conversion call is really inserted in this case.
And, if the function is called with an argument of type integer, the parser will try to convert that to text:

```sql
SELECT substr(1234, 3);
```

**ERROR:** function substr(integer, integer) does not exist

**HINT:** No function matches the given name and argument types. You might need to add explicit type casts.

This does not work because integer does not have an implicit cast to text. An explicit cast will work, however:

```sql
SELECT substr(CAST (1234 AS text), 3);
```

```sql
<table>
<thead>
<tr>
<th>substr</th>
</tr>
</thead>
<tbody>
<tr>
<td>34</td>
</tr>
</tbody>
</table>
```

(1 row)

## 10.4. Value Storage

Values to be inserted into a table are converted to the destination column’s data type according to the following steps.

**Value Storage Type Conversion**

1. Check for an exact match with the target.
2. Otherwise, try to convert the expression to the target type. This is possible if an assignment cast between the two types is registered in the `pg_cast` catalog (see CREATE CAST). Alternatively, if the expression is an unknown-type literal, the contents of the literal string will be fed to the input conversion routine for the target type.
3. Check to see if there is a sizing cast for the target type. A sizing cast is a cast from that type to itself. If one is found in the `pg_cast` catalog, apply it to the expression before storing into the destination column. The implementation function for such a cast always takes an extra parameter of type integer, which receives the destination column’s `atttypmod` value (typically its declared length, although the interpretation of `atttypmod` varies for different data types), and it may take a third boolean parameter that says whether the cast is explicit or implicit. The cast function is responsible for applying any length-dependent semantics such as size checking or truncation.

**Example 10-9. character Storage Type Conversion**

For a target column declared as `character(20)` the following statement shows that the stored value is sized correctly:

```sql
CREATE TABLE vv (v character(20));
INSERT INTO vv SELECT 'abc' || 'def';
SELECT v, octet_length(v) FROM vv;
```

```sql
<table>
<thead>
<tr>
<th>v</th>
<th>octet_length</th>
</tr>
</thead>
<tbody>
<tr>
<td>abcdef</td>
<td>20</td>
</tr>
</tbody>
</table>
```

(1 row)
Chapter 10. Type Conversion

What has really happened here is that the two unknown literals are resolved to `text` by default, allowing the `||` operator to be resolved as `text` concatenation. Then the `text` result of the operator is converted to `bpchar` ("blank-padded char", the internal name of the `character` data type) to match the target column type. (Since the conversion from `text` to `bpchar` is binary-coercible, this conversion does not insert any real function call.) Finally, the sizing function `bpchar(bpchar, integer, boolean)` is found in the system catalog and applied to the operator’s result and the stored column length. This type-specific function performs the required length check and addition of padding spaces.

10.5. `UNION`, `CASE`, and Related Constructs

SQL `UNION` constructs must match up possibly dissimilar types to become a single result set. The resolution algorithm is applied separately to each output column of a union query. The `INTERSECT` and `EXCEPT` constructs resolve dissimilar types in the same way as `UNION`. The `CASE`, `ARRAY`, `VALUES`, `GREATEST` and `LEAST` constructs use the identical algorithm to match up their component expressions and select a result data type.

Type Resolution for `UNION`, `CASE`, and Related Constructs

1. If all inputs are of the same type, and it is not `unknown`, resolve as that type.
2. If any input is of a domain type, treat it as being of the domain’s base type for all subsequent steps.
3. If all inputs are of type `unknown`, resolve as type `text` (the preferred type of the string category). Otherwise, `unknown` inputs are ignored.
4. If the non-unknown inputs are not all of the same type category, fail.
5. Choose the first non-unknown input type which is a preferred type in that category, if there is one.
6. Otherwise, choose the last non-unknown input type that allows all the preceding non-unknown inputs to be implicitly converted to it. (There always is such a type, since at least the first type in the list must satisfy this condition.)
7. Convert all inputs to the selected type. Fail if there is not a conversion from a given input to the selected type.

Some examples follow.

Example 10-10. Type Resolution with Underspecified Types in a Union

```
SELECT text 'a' AS "text" UNION SELECT 'b';
```

```
text
---
a
b
```

4. Somewhat like the treatment of domain inputs for operators and functions, this behavior allows a domain type to be preserved through a `UNION` or similar construct, so long as the user is careful to ensure that all inputs are implicitly or explicitly of that exact type. Otherwise the domain’s base type will be preferred.
Here, the unknown-type literal ‘b’ will be resolved to type text.

Example 10-11. Type Resolution in a Simple Union

```sql
SELECT 1.2 AS "numeric" UNION SELECT 1;
```

<table>
<thead>
<tr>
<th>numeric</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
</tr>
<tr>
<td>1.2</td>
</tr>
</tbody>
</table>

(2 rows)
The literal 1.2 is of type numeric, and the integer value 1 can be cast implicitly to numeric, so that type is used.

Example 10-12. Type Resolution in a Transposed Union

```sql
SELECT 1 AS "real" UNION SELECT CAST('2.2' AS REAL);
```

<table>
<thead>
<tr>
<th>real</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
</tr>
<tr>
<td>2.2</td>
</tr>
</tbody>
</table>

(2 rows)
Here, since type real cannot be implicitly cast to integer, but integer can be implicitly cast to real, the union result type is resolved as real.

Example 10-13. Type Resolution in a Nested Union

```sql
SELECT NULL UNION SELECT NULL UNION SELECT 1;
```

ERROR: UNION types text and integer cannot be matched
This failure occurs because PostgreSQL treats multiple UNIONS as a nest of pairwise operations; that is, this input is the same as

```sql
(SELECT NULL UNION SELECT NULL) UNION SELECT 1;
```
The inner UNION is resolved as emitting type text, according to the rules given above. Then the outer UNION has inputs of types text and integer, leading to the observed error. The problem can be fixed by ensuring that the leftmost UNION has at least one input of the desired result type.

INTERSECT and EXCEPT operations are likewise resolved pairwise. However, the other constructs described in this section consider all of their inputs in one resolution step.
Chapter 11. Indexes

Indexes are a common way to enhance database performance. An index allows the database server to find and retrieve specific rows much faster than it could do without an index. But indexes also add overhead to the database system as a whole, so they should be used sensibly.

11.1. Introduction

Suppose we have a table similar to this:

```sql
CREATE TABLE test1 (
    id integer,
    content varchar
);
```

and the application issues many queries of the form:

```sql
SELECT content FROM test1 WHERE id = constant;
```

With no advance preparation, the system would have to scan the entire `test1` table, row by row, to find all matching entries. If there are many rows in `test1` and only a few rows (perhaps zero or one) that would be returned by such a query, this is clearly an inefficient method. But if the system has been instructed to maintain an index on the `id` column, it can use a more efficient method for locating matching rows. For instance, it might only have to walk a few levels deep into a search tree.

A similar approach is used in most non-fiction books: terms and concepts that are frequently looked up by readers are collected in an alphabetic index at the end of the book. The interested reader can scan the index relatively quickly and flip to the appropriate page(s), rather than having to read the entire book to find the material of interest. Just as it is the task of the author to anticipate the items that readers are likely to look up, it is the task of the database programmer to foresee which indexes will be useful.

The following command can be used to create an index on the `id` column, as discussed:

```sql
CREATE INDEX test1_id_index ON test1 (id);
```

The name `test1_id_index` can be chosen freely, but you should pick something that enables you to remember later what the index was for.

To remove an index, use the `DROP INDEX` command. Indexes can be added to and removed from tables at any time.

Once an index is created, no further intervention is required: the system will update the index when the table is modified, and it will use the index in queries when it thinks doing so would be more efficient than a sequential table scan. But you might have to run the `ANALYZE` command regularly to update statistics to allow the query planner to make educated decisions. See Chapter 14 for information about how to find out whether an index is used and when and why the planner might choose not to use an index.

Indexes can also benefit `UPDATE` and `DELETE` commands with search conditions. Indexes can moreover be used in join searches. Thus, an index defined on a column that is part of a join condition can also significantly speed up queries with joins.

Creating an index on a large table can take a long time. By default, PostgreSQL allows reads (`SELECT` statements) to occur on the table in parallel with index creation, but writes (`INSERT`, `UPDATE`, `DELETE`) to the table are serialized. This is because it is not possible to concurrently update the table and the index. If you need to create an index on a table with a large number of rows, you might want to increase the number of parallel reads allowed by PostgreSQL.
UPDATE, DELETE) are blocked until the index build is finished. In production environments this is often unacceptable. It is possible to allow writes to occur in parallel with index creation, but there are several caveats to be aware of — for more information see Building Indexes Concurrently.

After an index is created, the system has to keep it synchronized with the table. This adds overhead to data manipulation operations. Therefore indexes that are seldom or never used in queries should be removed.

11.2. Index Types

PostgreSQL provides several index types: B-tree, Hash, GiST, SP-GiST, GIN and BRIN. Each index type uses a different algorithm that is best suited to different types of queries. By default, the CREATE INDEX command creates B-tree indexes, which fit the most common situations.

B-trees can handle equality and range queries on data that can be sorted into some ordering. In particular, the PostgreSQL query planner will consider using a B-tree index whenever an indexed column is involved in a comparison using one of these operators:

<  <=  =  >=  >

Constructs equivalent to combinations of these operators, such as BETWEEN and IN, can also be implemented with a B-tree index search. Also, an IS NULL or IS NOT NULL condition on an index column can be used with a B-tree index.

The optimizer can also use a B-tree index for queries involving the pattern matching operators LIKE and ~ if the pattern is a constant and is anchored to the beginning of the string — for example, col LIKE ‘foo%’ or col ~ '^foo', but not col LIKE '%bar'. However, if your database does not use the C locale you will need to create the index with a special operator class to support indexing of pattern-matching queries; see Section 11.9 below. It is also possible to use B-tree indexes for ILIKE and ~*, but only if the pattern starts with non-alphabetic characters, i.e., characters that are not affected by upper/lower case conversion.

B-tree indexes can also be used to retrieve data in sorted order. This is not always faster than a simple scan and sort, but it is often helpful.

Hash indexes can only handle simple equality comparisons. The query planner will consider using a hash index whenever an indexed column is involved in a comparison using the = operator. The following command is used to create a hash index:

```
CREATE INDEX name ON table USING HASH {column};
```

Caution

Hash index operations are not presently WAL-logged, so hash indexes might need to be rebuilt with REINDEX after a database crash if there were unwritten changes. Also, changes to hash indexes are not replicated over streaming or file-based replication after the initial base backup, so they give wrong answers to queries that subsequently use them. For these reasons, hash index use is presently discouraged.
GiST indexes are not a single kind of index, but rather an infrastructure within which many different indexing strategies can be implemented. Accordingly, the particular operators with which a GiST index can be used vary depending on the indexing strategy (the operator class). As an example, the standard distribution of PostgreSQL includes GiST operator classes for several two-dimensional geometric data types, which support indexed queries using these operators:

```
<< & << & > & >> << | & < | | & > | > @ =~ @ ~ =~
```

(See Section 9.11 for the meaning of these operators.) The GiST operator classes included in the standard distribution are documented in Table 61-1. Many other GiST operator classes are available in the contrib collection or as separate projects. For more information see Chapter 61.

GiST indexes are also capable of optimizing “nearest-neighbor” searches, such as

```
SELECT * FROM places ORDER BY location <-> point '(101,456)' LIMIT 10;
```

which finds the ten places closest to a given target point. The ability to do this is again dependent on the particular operator class being used. In Table 61-1, operators that can be used in this way are listed in the column “Ordering Operators”.

SP-GiST indexes, like GiST indexes, offer an infrastructure that supports various kinds of searches. SP-GiST permits implementation of a wide range of different non-balanced disk-based data structures, such as quadtrees, k-d trees, and radix trees (tries). As an example, the standard distribution of PostgreSQL includes SP-GiST operator classes for two-dimensional points, which support indexed queries using these operators:

```
<< > > =~ < @ ^ ^ >
```

(See Section 9.11 for the meaning of these operators.) The SP-GiST operator classes included in the standard distribution are documented in Table 62-1. For more information see Chapter 62.

GIN indexes are “inverted indexes” which are appropriate for data values that contain multiple component values, such as arrays. An inverted index contains a separate entry for each component value, and can efficiently handle queries that test for the presence of specific component values.

Like GiST and SP-GiST, GIN can support many different user-defined indexing strategies, and the particular operators with which a GIN index can be used vary depending on the indexing strategy. As an example, the standard distribution of PostgreSQL includes GIN operator classes for one-dimensional arrays, which support indexed queries using these operators:
Chapter 11. Indexes

<@
@
@
>&
&&
<br>
(See Section 9.18 for the meaning of these operators.) The GIN operator classes included in the standard distribution are documented in Table 63-1. Many other GIN operator classes are available in the contrib collection or as separate projects. For more information see Chapter 63.

BRIN indexes (a shorthand for Block Range INdexes) store summaries about the values stored in consecutive physical block ranges of a table. Like GiST, SP-GiST and GIN, BRIN can support many different indexing strategies, and the particular operators with which a BRIN index can be used vary depending on the indexing strategy. For data types that have a linear sort order, the indexed data corresponds to the minimum and maximum values of the values in the column for each block range. This supports indexed queries using these operators:

<
<=
-
>=
>

The BRIN operator classes included in the standard distribution are documented in Table 64-1. For more information see Chapter 64.

11.3. Multicolumn Indexes

An index can be defined on more than one column of a table. For example, if you have a table of this form:

```sql
CREATE TABLE test2 (  
   major int,  
   minor int,  
   name varchar  
);  
```

(say, you keep your /dev directory in a database...) and you frequently issue queries like:

```sql
SELECT name FROM test2 WHERE major = constant AND minor = constant;  
```

then it might be appropriate to define an index on the columns major and minor together, e.g.:

```sql
CREATE INDEX test2_mm_idx ON test2 (major, minor);  
```

Currently, only the B-tree, GiST, GIN, and BRIN index types support multicolumn indexes. Up to 32 columns can be specified. (This limit can be altered when building PostgreSQL; see the file pg_config_manual.h.)

A multicolumn B-tree index can be used with query conditions that involve any subset of the index’s columns, but the index is most efficient when there are constraints on the leading (leftmost) columns. The exact rule is that equality constraints on leading columns, plus any inequality constraints on the first column that does not have an equality constraint, will be used to limit the portion of the index that is scanned. Constraints on columns to the right of these columns are checked in the index, so they save
visits to the table proper, but they do not reduce the portion of the index that has to be scanned. For example, given an index on \((a, b, c)\) and a query condition \(\text{WHERE } a = 5 \text{ AND } b >= 42 \text{ AND } c < 77\), the index would have to be scanned from the first entry with \(a = 5\) and \(b = 42\) up through the last entry with \(a = 5\). Index entries with \(c >= 77\) would be skipped, but they’d still have to be scanned through. This index could in principle be used for queries that have constraints on \(b\) and/or \(c\) with no constraint on \(a\) — but the entire index would have to be scanned, so in most cases the planner would prefer a sequential table scan over using the index.

A multicolumn GiST index can be used with query conditions that involve any subset of the index’s columns. Conditions on additional columns restrict the entries returned by the index, but the condition on the first column is the most important one for determining how much of the index needs to be scanned. A GiST index will be relatively ineffective if its first column has only a few distinct values, even if there are many distinct values in additional columns.

A multicolumn GIN index can be used with query conditions that involve any subset of the index’s columns. Unlike B-tree or GiST, index search effectiveness is the same regardless of which index column(s) the query conditions use.

A multicolumn BRIN index can be used with query conditions that involve any subset of the index’s columns. Like GIN and unlike B-tree or GiST, index search effectiveness is the same regardless of which index column(s) the query conditions use. The only reason to have multiple BRIN indexes instead of one multicolumn BRIN index on a single table is to have a different `pages_per_range` storage parameter.

Of course, each column must be used with operators appropriate to the index type; clauses that involve other operators will not be considered.

Multicolumn indexes should be used sparingly. In most situations, an index on a single column is sufficient and saves space and time. Indexes with more than three columns are unlikely to be helpful unless the usage of the table is extremely stylized. See also Section 11.5 and Section 11.11 for some discussion of the merits of different index configurations.

### 11.4. Indexes and ORDER BY

In addition to simply finding the rows to be returned by a query, an index may be able to deliver them in a specific sorted order. This allows a query’s `ORDER BY` specification to be honored without a separate sorting step. Of the index types currently supported by PostgreSQL, only B-tree can produce sorted output — the other index types return matching rows in an unspecified, implementation-dependent order.

The planner will consider satisfying an `ORDER BY` specification either by scanning an available index that matches the specification, or by scanning the table in physical order and doing an explicit sort. For a query that requires scanning a large fraction of the table, an explicit sort is likely to be faster than using an index because it requires less disk I/O due to following a sequential access pattern. Indexes are more useful when only a few rows need be fetched. An important special case is `ORDER BY` in combination with `LIMIT n`: an explicit sort will have to process all the data to identify the first \(n\) rows, but if there is an index matching the `ORDER BY`, the first \(n\) rows can be retrieved directly, without scanning the remainder at all.

By default, B-tree indexes store their entries in ascending order with nulls last. This means that a forward scan of an index on column \(x\) produces output satisfying `ORDER BY x` (or more verbosely, `ORDER BY x ASC NULLS LAST`). The index can also be scanned backward, producing output satisfying `ORDER BY x DESC` (or more verbosely, `ORDER BY x DESC NULLS FIRST`, since `NULLS FIRST` is the default for `ORDER BY DESC`).
You can adjust the ordering of a B-tree index by including the options `ASC`, `DESC`, `NULLS FIRST`, and/or `NULLS LAST` when creating the index; for example:

```
CREATE INDEX test2_info_nulls_low ON test2 (info NULLS FIRST);
CREATE INDEX test3_desc_index ON test3 (id DESC NULLS LAST);
```

An index stored in ascending order with nulls first can satisfy either `ORDER BY x ASC NULLS FIRST` or `ORDER BY x DESC NULLS LAST` depending on which direction it is scanned in.

You might wonder why bother providing all four options, when two options together with the possibility of backward scan would cover all the variants of `ORDER BY`. In single-column indexes the options are indeed redundant, but in multicoloumn indexes they can be useful. Consider a two-column index on `(x, y)`: this can satisfy `ORDER BY x, y` if we scan forward, or `ORDER BY x DESC, y DESC` if we scan backward. But it might be that the application frequently needs to use `ORDER BY x ASC, y DESC`. There is no way to get that ordering from a plain index, but it is possible if the index is defined as `(x ASC, y DESC)` or `(x DESC, y ASC).

Obviously, indexes with non-default sort orderings are a fairly specialized feature, but sometimes they can produce tremendous speedups for certain queries. Whether it’s worth maintaining such an index depends on how often you use queries that require a special sort ordering.

11.5. Combining Multiple Indexes

A single index scan can only use query clauses that use the index’s columns with operators of its operator class and are joined with `AND`. For example, given an index on `(a, b)` a query condition like `WHERE a = 5 AND b = 6` could use the index, but a query like `WHERE a = 5 OR b = 6` could not directly use the index.

Fortunately, PostgreSQL has the ability to combine multiple indexes (including multiple uses of the same index) to handle cases that cannot be implemented by single index scans. The system can form `AND` and `OR` conditions across several index scans. For example, a query like `WHERE x = 42 OR x = 47 OR x = 53 OR x = 99` could be broken down into four separate scans of an index on `x`, each scan using one of the query clauses. The results of these scans are then ORed together to produce the result. Another example is that if we have separate indexes on `x` and `y`, one possible implementation of a query like `WHERE x = 5 AND y = 6` is to use each index with the appropriate query clause and then AND together the index results to identify the result rows.

To combine multiple indexes, the system scans each needed index and prepares a `bitmap` in memory giving the locations of table rows that are reported as matching that index’s conditions. The bitmaps are then ANDed and ORed together as needed by the query. Finally, the actual table rows are visited and returned. The table rows are visited in physical order, because that is how the bitmap is laid out; this means that any ordering of the original indexes is lost, and so a separate sort step will be needed if the query has an `ORDER BY` clause. For this reason, and because each additional index scan adds extra time, the planner will sometimes choose to use a simple index scan even though additional indexes are available that could have been used as well.

In all but the simplest applications, there are various combinations of indexes that might be useful, and the database developer must make trade-offs to decide which indexes to provide. Sometimes multicoloumn indexes are best, but sometimes it’s better to create separate indexes and rely on the index-combination feature. For example, if your workload includes a mix of queries that sometimes involve only column `x`, sometimes only column `y`, and sometimes both columns, you might choose to create two separate indexes on `x` and `y`, relying on index combination to process the queries that use both columns. You could also create a multicoloumn index on `(x, y)`. This index would typically be more efficient than index combination for queries involving both columns, but as discussed in Section
11.3. it would be almost useless for queries involving only \( y \), so it should not be the only index. A combination of the multicolumn index and a separate index on \( y \) would serve reasonably well. For queries involving only \( x \), the multicolumn index could be used, though it would be larger and hence slower than an index on \( x \) alone. The last alternative is to create all three indexes, but this is probably only reasonable if the table is searched much more often than it is updated and all three types of query are common. If one of the types of query is much less common than the others, you’d probably settle for creating just the two indexes that best match the common types.

### 11.6. Unique Indexes

Indexes can also be used to enforce uniqueness of a column’s value, or the uniqueness of the combined values of more than one column.

```sql
CREATE UNIQUE INDEX name ON table (column [, ...]);
```

Currently, only B-tree indexes can be declared unique.

When an index is declared unique, multiple table rows with equal indexed values are not allowed. Null values are not considered equal. A multicolumn unique index will only reject cases where all indexed columns are equal in multiple rows.

PostgreSQL automatically creates a unique index when a unique constraint or primary key is defined for a table. The index covers the columns that make up the primary key or unique constraint (a multicolumn index, if appropriate), and is the mechanism that enforces the constraint.

**Note:** There’s no need to manually create indexes on unique columns; doing so would just duplicate the automatically-created index.

### 11.7. Indexes on Expressions

An index column need not be just a column of the underlying table, but can be a function or scalar expression computed from one or more columns of the table. This feature is useful to obtain fast access to tables based on the results of computations.

For example, a common way to do case-insensitive comparisons is to use the `lower` function:

```sql
SELECT * FROM test1 WHERE lower(col1) = 'value';
```

This query can use an index if one has been defined on the result of the `lower(col1)` function:

```sql
CREATE INDEX test1_lower_col1_idx ON test1 (lower(col1));
```

Expression indexes also allow control over the scope of unique indexes. For example, this unique index prevents duplicate integer values from being stored in a `double precision`-typed column:

```sql
CREATE UNIQUE INDEX test1_uniq_int ON tests ((floor(double_col)));
```
If we were to declare this index UNIQue, it would prevent creation of rows whose col1 values differ only in case, as well as rows whose col1 values are actually identical. Thus, indexes on expressions can be used to enforce constraints that are not definable as simple unique constraints.

As another example, if one often does queries like:

```sql
SELECT * FROM people WHERE (first_name || ' ' || last_name) = 'John Smith';
```

then it might be worth creating an index like this:

```sql
CREATE INDEX people_names ON people ((first_name || ' ' || last_name));
```

The syntax of the `CREATE INDEX` command normally requires writing parentheses around index expressions, as shown in the second example. The parentheses can be omitted when the expression is just a function call, as in the first example.

Index expressions are relatively expensive to maintain, because the derived expression(s) must be computed for each row upon insertion and whenever it is updated. However, the index expressions are not recomputed during an indexed search, since they are already stored in the index. In both examples above, the system sees the query as just `WHERE indexedcolumn = 'constant'` and so the speed of the search is equivalent to any other simple index query. Thus, indexes on expressions are useful when retrieval speed is more important than insertion and update speed.

### 11.8. Partial Indexes

A partial index is an index built over a subset of a table; the subset is defined by a conditional expression (called the predicate of the partial index). The index contains entries only for those table rows that satisfy the predicate. Partial indexes are a specialized feature, but there are several situations in which they are useful.

One major reason for using a partial index is to avoid indexing common values. Since a query searching for a common value (one that accounts for more than a few percent of all the table rows) will not use the index anyway, there is no point in keeping those rows in the index at all. This reduces the size of the index, which will speed up those queries that do use the index. It will also speed up many table update operations because the index does not need to be updated in all cases. Example 11-1 shows a possible application of this idea.

#### Example 11-1. Setting up a Partial Index to Exclude Common Values

Suppose you are storing web server access logs in a database. Most accesses originate from the IP address range of your organization but some are from elsewhere (say, employees on dial-up connections). If your searches by IP are primarily for outside accesses, you probably do not need to index the IP range that corresponds to your organization’s subnet.

Assume a table like this:

```sql
CREATE TABLE access_log (
    url varchar,
    client_ip inet,
    ...
);
```

To create a partial index that suits our example, use a command such as this:

```sql
CREATE INDEX access_log_client_ip_ix ON access_log (client_ip)
```
WHERE NOT (client_ip > inet '192.168.100.0' AND client_ip < inet '192.168.100.255');

A typical query that can use this index would be:

```
SELECT *
FROM access_log
WHERE url = '/index.html' AND client_ip = inet '212.78.10.32';
```

A query that cannot use this index is:

```
SELECT *
FROM access_log
WHERE client_ip = inet '192.168.100.23';
```

Observe that this kind of partial index requires that the common values be predetermined, so such partial indexes are best used for data distributions that do not change. The indexes can be recreated occasionally to adjust for new data distributions, but this adds maintenance effort.

Another possible use for a partial index is to exclude values from the index that the typical query workload is not interested in; this is shown in Example 11-2. This results in the same advantages as listed above, but it prevents the “uninteresting” values from being accessed via that index, even if an index scan might be profitable in that case. Obviously, setting up partial indexes for this kind of scenario will require a lot of care and experimentation.

**Example 11-2. Setting up a Partial Index to Exclude Uninteresting Values**

If you have a table that contains both billed and unbilled orders, where the unbilled orders take up a small fraction of the total table and yet those are the most-accessed rows, you can improve performance by creating an index on just the unbilled rows. The command to create the index would look like this:

```
CREATE INDEX orders_unbilled_index ON orders (order_nr)
WHERE billed is not true;
```

A possible query to use this index would be:

```
SELECT * FROM orders WHERE billed is not true AND order_nr < 10000;
```

However, the index can also be used in queries that do not involve `order_nr` at all, e.g.:

```
SELECT * FROM orders WHERE billed is not true AND amount > 5000.00;
```

This is not as efficient as a partial index on the `amount` column would be, since the system has to scan the entire index. Yet, if there are relatively few unbilled orders, using this partial index just to find the unbilled orders could be a win.

Note that this query cannot use this index:

```
SELECT * FROM orders WHERE order_nr = 3501;
```

The order 3501 might be among the billed or unbilled orders.

Example 11-2 also illustrates that the indexed column and the column used in the predicate do not need to match. PostgreSQL supports partial indexes with arbitrary predicates, so long as only columns of the table being indexed are involved. However, keep in mind that the predicate must match the conditions used in the queries that are supposed to benefit from the index. To be precise, a partial index can be used in a query only if the system can recognize that the `WHERE` condition of the query mathematically implies the predicate of the index. PostgreSQL does not have a sophisticated theorem prover that can recognize mathematically equivalent expressions that are written in different forms. (Not only is such a general theorem prover extremely difficult to create, it would probably be too
slow to be of any real use.) The system can recognize simple inequality implications, for example
"x < 1" implies "x < 2"; otherwise the predicate condition must exactly match part of the query’s
WHERE condition or the index will not be recognized as usable. Matching takes place at query planning
time, not at run time. As a result, parameterized query clauses do not work with a partial index. For
example a prepared query with a parameter might specify "x < ?" which will never imply "x < 2" for
all possible values of the parameter.

A third possible use for partial indexes does not require the index to be used in queries at all. The idea
here is to create a unique index over a subset of a table, as in Example 11-3. This enforces uniqueness
among the rows that satisfy the index predicate, without constraining those that do not.

Example 11-3. Setting up a Partial Unique Index

Suppose that we have a table describing test outcomes. We wish to ensure that there is only one
“successful” entry for a given subject and target combination, but there might be any number of
“unsuccessful” entries. Here is one way to do it:

CREATE TABLE tests (
  subject text,
  target text,
  success boolean,
  ...
);
CREATE UNIQUE INDEX tests_success_constraint ON tests (subject, target)
  WHERE success;

This is a particularly efficient approach when there are few successful tests and many unsuccessful
ones.

This index allows only one null in the indexed column by using a partial index clause to process only
null column values, and using an expression index clause to index true instead of null:

CREATE UNIQUE INDEX tests_target_one_null ON tests ((target IS NULL)) WHERE target IS NULL;

Finally, a partial index can also be used to override the system’s query plan choices. Also, data sets
with peculiar distributions might cause the system to use an index when it really should not. In that
case the index can be set up so that it is not available for the offending query. Normally, PostgreSQL
makes reasonable choices about index usage (e.g., it avoids them when retrieving common values, so
the earlier example really only saves index size, it is not required to avoid index usage), and grossly
incorrect plan choices are cause for a bug report.

Keep in mind that setting up a partial index indicates that you know at least as much as the query
planner knows, in particular you know when an index might be profitable. Forming this knowledge
requires experience and understanding of how indexes in PostgreSQL work. In most cases, the advan-
tage of a partial index over a regular index will be minimal.

More information about partial indexes can be found in The case for partial indexes, Partial indexing
in POSTGRES: research project, and Generalized Partial Indexes (cached version).

11.9. Operator Classes and Operator Families

An index definition can specify an operator class for each column of an index.

CREATE INDEX name ON table (column opclass [sort options] [, ...]);

Chapter 11. Indexes
The operator class identifies the operators to be used by the index for that column. For example, a B-tree index on the type \texttt{int4} would use the \texttt{int4_ops} class; this operator class includes comparison functions for values of type \texttt{int4}. In practice the default operator class for the column’s data type is usually sufficient. The main reason for having operator classes is that for some data types, there could be more than one meaningful index behavior. For example, we might want to sort a complex-number data type either by absolute value or by real part. We could do this by defining two operator classes for the data type and then selecting the proper class when making an index. The operator class determines the basic sort ordering (which can then be modified by adding sort options \texttt{COLLATE}, \texttt{ASC}/\texttt{DESC} and/or \texttt{NULLS FIRST}/\texttt{NULLS LAST}).

There are also some built-in operator classes besides the default ones:

- The operator classes \texttt{text_pattern_ops}, \texttt{varchar_pattern_ops}, and \texttt{bpchar_pattern_ops} support B-tree indexes on the types \texttt{text}, \texttt{varchar}, and \texttt{char} respectively. The difference from the default operator classes is that the values are compared strictly character by character rather than according to the locale-specific collation rules. This makes these operator classes suitable for use by queries involving pattern matching expressions (\texttt{LIKE} or POSIX regular expressions) when the database does not use the standard “C” locale. As an example, you might index a \texttt{varchar} column like this:

```
CREATE INDEX test_index ON test_table (col varchar_pattern_ops);
```

Note that you should also create an index with the default operator class if you want queries involving ordinary \texttt{<}, \texttt{<=}, \texttt{>,} or \texttt{>=} comparisons to use an index. Such queries cannot use the \texttt{xxx_pattern_ops} operator classes. (Ordinary equality comparisons can use these operator classes, however.) It is possible to create multiple indexes on the same column with different operator classes. If you do use the C locale, you do not need the \texttt{xxx_pattern_ops} operator classes, because an index with the default operator class is usable for pattern-matching queries in the C locale.

The following query shows all defined operator classes:

```
SELECT am.amname AS index_method, 
opc.opcname AS opclass_name, 
opc.opcintype::regtype AS indexed_type, 
opc.opcdefault AS is_default 
FROM pg_am am, pg_opclass opc 
WHERE opc.opcmethod = am.oid 
ORDER BY index_method, opclass_name;
```

An operator class is actually just a subset of a larger structure called an \textit{operator family}. In cases where several data types have similar behaviors, it is frequently useful to define cross-data-type operators and allow these to work with indexes. To do this, the operator classes for each of the types must be grouped into the same operator family. The cross-type operators are members of the family, but are not associated with any single class within the family.

This expanded version of the previous query shows the operator family each operator class belongs to:

```
SELECT am.amname AS index_method, 
opc.opcname AS opclass_name, 
opf.opfname AS opfamily_name, 
opc.opcintype::regtype AS indexed_type, 
opc.opcdefault AS is_default 
FROM pg_am am, pg_opclass opc, pg_opfamily opf 
WHERE opc.opcmethod = am.oid 
    AND opc.opcfamily = opf.oid 
ORDER BY index_method, opclass_name, opfamily_name;
```
Chapter 11. Indexes

FROM pg_am am, pg_opclass opc, pg_opfamily opf
WHERE opc.opcmethod = am.oid AND
   opc.opcfamily = opf.oid
ORDER BY index_method, opclass_name;

This query shows all defined operator families and all the operators included in each family:

SELECT am.amname AS index_method,
       opf.opfname AS opfamily_name,
       amop.amopopr::regoperator AS opfamily_operator
FROM pg_am am, pg_opfamily opf, pg_amop amop
WHERE opf.opfmethod = am.oid AND
       amop.amopfamily = opf.oid
ORDER BY index_method, opfamily_name, opfamily_operator;

11.10. Indexes and Collations

An index can support only one collation per index column. If multiple collations are of interest, multiple indexes may be needed.

Consider these statements:

CREATE TABLE test1c (id integer,
                      content varchar COLLATE "x")
CREATE INDEX test1c_content_index ON test1c (content);

The index automatically uses the collation of the underlying column. So a query of the form

SELECT * FROM test1c WHERE content > constant;

could use the index, because the comparison will by default use the collation of the column. However, this index cannot accelerate queries that involve some other collation. So if queries of the form, say,

SELECT * FROM test1c WHERE content > constant COLLATE "y";

are also of interest, an additional index could be created that supports the "y" collation, like this:

CREATE INDEX test1c_content_y_index ON test1c (content COLLATE "y");

11.11. Index-Only Scans

All indexes in PostgreSQL are secondary indexes, meaning that each index is stored separately from the table’s main data area (which is called the table’s heap in PostgreSQL terminology). This means that in an ordinary index scan, each row retrieval requires fetching data from both the index and the heap. Furthermore, while the index entries that match a given indexable WHERE condition are usually
close together in the index, the table rows they reference might be anywhere in the heap. The heap-access portion of an index scan thus involves a lot of random access into the heap, which can be slow, particularly on traditional rotating media. (As described in Section 11.5, bitmap scans try to alleviate this cost by doing the heap accesses in sorted order, but that only goes so far.)

To solve this performance problem, PostgreSQL supports index-only scans, which can answer queries from an index alone without any heap access. The basic idea is to return values directly out of each index entry instead of consulting the associated heap entry. There are two fundamental restrictions on when this method can be used:

1. The index type must support index-only scans. B-tree indexes always do. GiST and SP-GiST indexes support index-only scans for some operator classes but not others. Other index types have no support. The underlying requirement is that the index must physically store, or else be able to reconstruct, the original data value for each index entry. As a counterexample, GIN indexes cannot support index-only scans because each index entry typically holds only part of the original data value.

2. The query must reference only columns stored in the index. For example, given an index on columns x and y of a table that also has a column z, these queries could use index-only scans:

   ```sql
   SELECT x, y FROM tab WHERE x = 'key';
   SELECT x FROM tab WHERE x = 'key' AND y < 42;
   ```

   but these queries could not:

   ```sql
   SELECT x, z FROM tab WHERE x = 'key';
   SELECT x FROM tab WHERE x = 'key' AND z < 42;
   ```

   (Expression indexes and partial indexes complicate this rule, as discussed below.)

If these two fundamental requirements are met, then all the data values required by the query are available from the index, so an index-only scan is physically possible. But there is an additional requirement for any table scan in PostgreSQL: it must verify that each retrieved row be “visible” to the query’s MVCC snapshot, as discussed in Chapter 13. Visibility information is not stored in index entries, only in heap entries; so at first glance it would seem that every row retrieval would require a heap access anyway. And this is indeed the case, if the table row has been modified recently. However, for seldom-changing data there is a way around this problem. PostgreSQL tracks, for each page in a table’s heap, whether all rows stored in that page are old enough to be visible to all current and future transactions. This information is stored in a bit in the table’s visibility map. An index-only scan, after finding a candidate index entry, checks the visibility map bit for the corresponding heap page. If it’s set, the row is known visible and so the data can be returned with no further work. If it’s not set, the heap entry must be visited to find out whether it’s visible, so no performance advantage is gained over a standard index scan. Even in the successful case, this approach trades visibility map accesses for heap accesses; but since the visibility map is four orders of magnitude smaller than the heap it describes, far less physical I/O is needed to access it. In most situations the visibility map remains cached in memory all the time.

In short, while an index-only scan is possible given the two fundamental requirements, it will be a win only if a significant fraction of the table’s heap pages have their all-visible map bits set. But tables in which a large fraction of the rows are unchanging are common enough to make this type of scan very useful in practice.

To make effective use of the index-only scan feature, you might choose to create indexes in which only the leading columns are meant to match WHERE clauses, while the trailing columns hold “payload” data to be returned by a query. For example, if you commonly run queries like

```sql
SELECT y FROM tab WHERE x = 'key';
```
the traditional approach to speeding up such queries would be to create an index on \( x \) only. However, an index on \((x, y)\) would offer the possibility of implementing this query as an index-only scan. As previously discussed, such an index would be larger and hence more expensive than an index on \( x \) alone, so this is attractive only if the table is known to be mostly static. Note it’s important that the index be declared on \((x, y)\) not \((y, x)\), as for most index types (particularly B-trees) searches that do not constrain the leading index columns are not very efficient.

In principle, index-only scans can be used with expression indexes. For example, given an index on \( f(x) \) where \( x \) is a table column, it should be possible to execute

\[
\text{SELECT } f(x) \text{ FROM tab WHERE } f(x) < 1;
\]

as an index-only scan; and this is very attractive if \( f() \) is an expensive-to-compute function. However, PostgreSQL’s planner is currently not very smart about such cases. It considers a query to be potentially executable by index-only scan only when all \textit{columns} needed by the query are available from the index. In this example, \( x \) is not needed except in the context \( f(x) \), but the planner does not notice that and concludes that an index-only scan is not possible. If an index-only scan seems sufficiently worthwhile, this can be worked around by declaring the index to be on \((f(x), x)\), where the second column is not expected to be used in practice but is just there to convince the planner that an index-only scan is possible. An additional caveat, if the goal is to avoid recalculating \( f(x) \), is that the planner won’t necessarily match uses of \( f(x) \) that aren’t in indexable \textit{WHERE} clauses to the index column. It will usually get this right in simple queries such as shown above, but not in queries that involve joins. These deficiencies may be remedied in future versions of PostgreSQL.

Partial indexes also have interesting interactions with index-only scans. Consider the partial index shown in Example 11-3:

\[
\text{CREATE UNIQUE INDEX tests_success_constraint ON tests (subject, target) WHERE success;}
\]

In principle, we could do an index-only scan on this index to satisfy a query like

\[
\text{SELECT target FROM tests WHERE subject = 'some-subject' AND success;}
\]

But there’s a problem: the \textit{WHERE} clause refers to \textit{success} which is not available as a result column of the index. Nonetheless, an index-only scan is possible because the plan does not need to recheck that part of the \textit{WHERE} clause at runtime; all entries found in the index necessarily have \textit{success} = \textit{true} so this need not be explicitly checked in the plan. PostgreSQL versions 9.6 and later will recognize such cases and allow index-only scans to be generated, but older versions will not.

### 11.12. Examining Index Usage

Although indexes in PostgreSQL do not need maintenance or tuning, it is still important to check which indexes are actually used by the real-life query workload. Examining index usage for an individual query is done with the \texttt{EXPLAIN} command; its application for this purpose is illustrated in Section 14.1. It is also possible to gather overall statistics about index usage in a running server, as described in Section 28.2.

It is difficult to formulate a general procedure for determining which indexes to create. There are a number of typical cases that have been shown in the examples throughout the previous sections. A good deal of experimentation is often necessary. The rest of this section gives some tips for that:

- Always run \texttt{ANALYZE} first. This command collects statistics about the distribution of the values in the table. This information is required to estimate the number of rows returned by a query, which
Chapter 11. Indexes

is needed by the planner to assign realistic costs to each possible query plan. In absence of any real statistics, some default values are assumed, which are almost certain to be inaccurate. Examining an application’s index usage without having run \texttt{ANALYZE} is therefore a lost cause. See Section 24.1.3 and Section 24.1.6 for more information.

- Use real data for experimentation. Using test data for setting up indexes will tell you what indexes you need for the test data, but that is all.

  It is especially fatal to use very small test data sets. While selecting 1000 out of 100000 rows could be a candidate for an index, selecting 1 out of 100 rows will hardly be, because the 100 rows probably fit within a single disk page, and there is no plan that can beat sequentially fetching 1 disk page.

  Also be careful when making up test data, which is often unavoidable when the application is not yet in production. Values that are very similar, completely random, or inserted in sorted order will skew the statistics away from the distribution that real data would have.

- When indexes are not used, it can be useful for testing to force their use. There are run-time parameters that can turn off various plan types (see Section 19.7.1). For instance, turning off sequential scans (\texttt{enable_seqscan}) and nested-loop joins (\texttt{enable_nestloop}), which are the most basic plans, will force the system to use a different plan. If the system still chooses a sequential scan or nested-loop join then there is probably a more fundamental reason why the index is not being used; for example, the query condition does not match the index. (What kind of query can use what kind of index is explained in the previous sections.)

- If forcing index usage does use the index, then there are two possibilities: Either the system is right and using the index is indeed not appropriate, or the cost estimates of the query plans are not reflecting reality. So you should time your query with and without indexes. The \texttt{EXPLAIN ANALYZE} command can be useful here.

- If it turns out that the cost estimates are wrong, there are, again, two possibilities. The total cost is computed from the per-row costs of each plan node times the selectivity estimate of the plan node. The costs estimated for the plan nodes can be adjusted via run-time parameters (described in Section 19.7.2). An inaccurate selectivity estimate is due to insufficient statistics. It might be possible to improve this by tuning the statistics-gathering parameters (see \texttt{ALTER TABLE}).

  If you do not succeed in adjusting the costs to be more appropriate, then you might have to resort to forcing index usage explicitly. You might also want to contact the PostgreSQL developers to examine the issue.
Chapter 12. Full Text Search

12.1. Introduction

Full Text Searching (or just text search) provides the capability to identify natural-language documents that satisfy a query, and optionally to sort them by relevance to the query. The most common type of search is to find all documents containing given query terms and return them in order of their similarity to the query. Notions of query and similarity are very flexible and depend on the specific application. The simplest search considers query as a set of words and similarity as the frequency of query words in the document.

Textual search operators have existed in databases for years. PostgreSQL has ~, ~*, LIKE, and ILIKE operators for textual data types, but they lack many essential properties required by modern information systems:

- There is no linguistic support, even for English. Regular expressions are not sufficient because they cannot easily handle derived words, e.g., satisfies and satisfy. You might miss documents that contain satisfies, although you probably would like to find them when searching for satisfy. It is possible to use OR to search for multiple derived forms, but this is tedious and error-prone (some words can have several thousand derivatives).
- They provide no ordering (ranking) of search results, which makes them ineffective when thousands of matching documents are found.
- They tend to be slow because there is no index support, so they must process all documents for every search.

Full text indexing allows documents to be preprocessed and an index saved for later rapid searching. Preprocessing includes:

Parsing documents into tokens. It is useful to identify various classes of tokens, e.g., numbers, words, complex words, email addresses, so that they can be processed differently. In principle token classes depend on the specific application, but for most purposes it is adequate to use a predefined set of classes. PostgreSQL uses a parser to perform this step. A standard parser is provided, and custom parsers can be created for specific needs.

Converting tokens into lexemes. A lexeme is a string, just like a token, but it has been normalized so that different forms of the same word are made alike. For example, normalization almost always includes folding upper-case letters to lower-case, and often involves removal of suffixes (such as s or es in English). This allows searches to find variant forms of the same word, without tediously entering all the possible variants. Also, this step typically eliminates stop words, which are words that are so common that they are useless for searching. (In short, then, tokens are raw fragments of the document text, while lexemes are words that are believed useful for indexing and searching.) PostgreSQL uses dictionaries to perform this step. Various standard dictionaries are provided, and custom ones can be created for specific needs.

Storing preprocessed documents optimized for searching. For example, each document can be represented as a sorted array of normalized lexemes. Along with the lexemes it is often desirable to store positional information to use for proximity ranking, so that a document that contains a more “dense” region of query words is assigned a higher rank than one with scattered query words.

Dictionaries allow fine-grained control over how tokens are normalized. With appropriate dictionaries, you can:
• Define stop words that should not be indexed.
• Map synonyms to a single word using Ispell.
• Map phrases to a single word using a thesaurus.
• Map different variations of a word to a canonical form using an Ispell dictionary.
• Map different variations of a word to a canonical form using Snowball stemmer rules.

A data type \texttt{tsvector} is provided for storing preprocessed documents, along with a type \texttt{tsquery} for representing processed queries (Section 8.11). There are many functions and operators available for these data types (Section 9.13), the most important of which is the match operator @@, which we introduce in Section 12.1.2. Full text searches can be accelerated using indexes (Section 12.9).

12.1.1. What Is a Document?

A document is the unit of searching in a full text search system; for example, a magazine article or email message. The text search engine must be able to parse documents and store associations of lexemes (key words) with their parent document. Later, these associations are used to search for documents that contain query words.

For searches within PostgreSQL, a document is normally a textual field within a row of a database table, or possibly a combination (concatenation) of such fields, perhaps stored in several tables or obtained dynamically. In other words, a document can be constructed from different parts for indexing and it might not be stored anywhere as a whole. For example:

\begin{verbatim}
SELECT title || ' ' || author || ' ' || abstract || ' ' || body AS document
FROM messages
WHERE mid = 12;

SELECT m.title || ' ' || m.author || ' ' || m.abstract || ' ' || d.body AS document
FROM messages m, docs d
WHERE mid = did AND mid = 12;
\end{verbatim}

\textbf{Note:} Actually, in these example queries, \texttt{coalesce} should be used to prevent a single \texttt{NULL} attribute from causing a \texttt{NULL} result for the whole document.

Another possibility is to store the documents as simple text files in the file system. In this case, the database can be used to store the full text index and to execute searches, and some unique identifier can be used to retrieve the document from the file system. However, retrieving files from outside the database requires superuser permissions or special function support, so this is usually less convenient than keeping all the data inside PostgreSQL. Also, keeping everything inside the database allows easy access to document metadata to assist in indexing and display.

For text search purposes, each document must be reduced to the preprocessed \texttt{tsvector} format. Searching and ranking are performed entirely on the \texttt{tsvector} representation of a document — the original text need only be retrieved when the document has been selected for display to a user. We therefore often speak of the \texttt{tsvector} as being the document, but of course it is only a compact representation of the full document.
12.1.2. Basic Text Matching

Full text searching in PostgreSQL is based on the match operator @@, which returns true if a tsvector (document) matches a tsquery (query). It doesn’t matter which data type is written first:

```sql
SELECT 'a fat cat sat on a mat and ate a fat rat':tsvector @@ 'cat & rat':tsquery;
?column?
----------
t
```

```sql
SELECT 'fat & cow':tsquery @@ 'a fat cat sat on a mat and ate a fat rat':tsvector;
?column?
----------
f
```

As the above example suggests, a tsquery is not just raw text, any more than a tsvector is. A tsquery contains search terms, which must be already-normalized lexemes, and may combine multiple terms using AND, OR, NOT, and FOLLOWED BY operators. (For syntax details see Section 8.11.2.) There are functions to_tsquery, plainto_tsquery, and phraseto_tsquery that are helpful in converting user-written text into a proper tsquery, primarily by normalizing words appearing in the text. Similarly, to_tsvector is used to parse and normalize a document string. So in practice a text search match would look more like this:

```sql
SELECT to_tsvector('fat cats ate fat rats') @@ to_tsquery('fat & rat');
?column?
----------
t
```

Observe that this match would not succeed if written as

```sql
SELECT 'fat cats ate fat rats':tsvector @@ to_tsquery('fat & rat');
?column?
----------
f
```

since here no normalization of the word rats will occur. The elements of a tsvector are lexemes, which are assumed already normalized, so rats does not match rat.

The @@ operator also supports text input, allowing explicit conversion of a text string to tsvector or tsquery to be skipped in simple cases. The variants available are:

```sql
tsvectors @@ tsquery
tquery @@ tsvectors
text @@ tsquery
text @@ text
```

The first two of these we saw already. The form text @@ tsquery is equivalent to to_tsvectors(x) @@ y. The form text @@ text is equivalent to to_tsvector(x) @@ plainto_tsquery(y).

Within a tsquery, the & (AND) operator specifies that both its arguments must appear in the document to have a match. Similarly, the | (OR) operator specifies that at least one of its arguments must appear, while the ! (NOT) operator specifies that its argument must not appear in order to have a match. For example, the query fat & ! rat matches documents that contain fat but not rat.
Searching for phrases is possible with the help of the \(<-\) (FOLLOWED BY) \texttt{tsquery} operator, which matches only if its arguments have matches that are adjacent and in the given order. For example:

\begin{verbatim}
SELECT to_tsvector('fatal error') @@ to_tsquery('fatal <- error');
?column?
-------
t
SELECT to_tsvector('error is not fatal') @@ to_tsquery('fatal <- error');
?column?
-------
f
\end{verbatim}

There is a more general version of the FOLLOWED BY operator having the form \(<N>\), where \(N\) is an integer standing for the difference between the positions of the matching lexemes. \(<1>\) is the same as \(<-\), while \(<2>\) allows exactly one other lexeme to appear between the matches, and so on. The \texttt{phraseto_tsquery} function makes use of this operator to construct a \texttt{tsquery} that can match a multi-word phrase when some of the words are stop words. For example:

\begin{verbatim}
SELECT phraseto_tsquery('cats ate rats');
phraseto_tsquery
-------------------
'cat' <- 'ate' <- 'rat'

SELECT phraseto_tsquery('the cats ate the rats');
phraseto_tsquery
-------------------
'cat' <- 'ate' <-2 'rat'
\end{verbatim}

A special case that’s sometimes useful is that \(<0>\) can be used to require that two patterns match the same word.

Parentheses can be used to control nesting of the \texttt{tsquery} operators. Without parentheses, \(|\) binds least tightly, then \& then \(<-\), and \! most tightly.

It’s worth noticing that the AND/OR/NOT operators mean something subtly different when they are within the arguments of a FOLLOWED BY operator than when they are not, because within FOLLOWED BY the exact position of the match is significant. For example, normally \(!x\) matches only documents that do not contain \(x\) anywhere. But \(!x \ <-\ y\) matches \(y\) if it is not immediately after an \(x\); an occurrence of \(x\) elsewhere in the document does not prevent a match. Another example is that \(x \ & \ y\) normally only requires that \(x\) and \(y\) both appear somewhere in the document, but \((x \ & \ y) \ <-\ z\) requires \(x\) and \(y\) to match at the same place, immediately before a \(z\). Thus this query behaves differently from \(x \ <\ y\ & \ z\ <-\ z\), which will match a document containing two separate sequences \(x\ z\) and \(y\ z\). (This specific query is useless as written, since \(x\) and \(y\) could not match at the same place; but with more complex situations such as prefix-match patterns, a query of this form could be useful.)

12.1.3. Configurations

The above are all simple text search examples. As mentioned before, full text search functionality includes the ability to do many more things: skip indexing certain words (stop words), process synonyms, and use sophisticated parsing, e.g., parse based on more than just white space. This function-
Chapter 12. Full Text Search

Ality is controlled by text search configurations. PostgreSQL comes with predefined configurations for many languages, and you can easily create your own configurations. (psql’s \dF command shows all available configurations.)

During installation an appropriate configuration is selected and default_text_search_config is set accordingly in postgresql.conf. If you are using the same text search configuration for the entire cluster you can use the value in postgresql.conf. To use different configurations throughout the cluster but the same configuration within any one database, use ALTER DATABASE ... SET. Otherwise, you can set default_text_search_config in each session.

Each text search function that depends on a configuration has an optional regconfig argument, so that the configuration to use can be specified explicitly. default_text_search_config is used only when this argument is omitted.

To make it easier to build custom text search configurations, a configuration is built up from simpler database objects. PostgreSQL’s text search facility provides four types of configuration-related database objects:

- **Text search parsers** break documents into tokens and classify each token (for example, as words or numbers).
- **Text search dictionaries** convert tokens to normalized form and reject stop words.
- **Text search templates** provide the functions underlying dictionaries. (A dictionary simply specifies a template and a set of parameters for the template.)
- **Text search configurations** select a parser and a set of dictionaries to use to normalize the tokens produced by the parser.

Text search parsers and templates are built from low-level C functions; therefore it requires C programming ability to develop new ones, and superuser privileges to install one into a database. (There are examples of add-on parsers and templates in the contrib/ area of the PostgreSQL distribution.) Since dictionaries and configurations just parameterize and connect together some underlying parsers and templates, no special privilege is needed to create a new dictionary or configuration. Examples of creating custom dictionaries and configurations appear later in this chapter.

12.2. Tables and Indexes

The examples in the previous section illustrated full text matching using simple constant strings. This section shows how to search table data, optionally using indexes.

12.2.1. Searching a Table

It is possible to do a full text search without an index. A simple query to print the title of each row that contains the word friend in its body field is:

```sql
SELECT title
FROM pgweb
WHERE to_tsvector('english', body) @@ to_tsquery('english', 'friend');
```

This will also find related words such as friends and friendly, since all these are reduced to the same normalized lexeme.

The query above specifies that the english configuration is to be used to parse and normalize the strings. Alternatively we could omit the configuration parameters:
SELECT title
FROM pgweb
WHERE to_tsvector(body) @@ to_tsquery('friend');

This query will use the configuration set by default_text_search_config.

A more complex example is to select the ten most recent documents that contain create and table in the title or body:

SELECT title
FROM pgweb
WHERE to_tsvector(title || ' ' || body) @@ to_tsquery('create & table')
ORDER BY last_mod_date DESC
LIMIT 10;

For clarity we omitted the coalesce function calls which would be needed to find rows that contain NULL in one of the two fields.

Although these queries will work without an index, most applications will find this approach too slow, except perhaps for occasional ad-hoc searches. Practical use of text searching usually requires creating an index.

12.2.2. Creating Indexes

We can create a GIN index (Section 12.9) to speed up text searches:

CREATE INDEX pgweb_idx ON pgweb USING GIN (to_tsvector('english', body));

Notice that the 2-argument version of to_tsvector is used. Only text search functions that specify a configuration name can be used in expression indexes (Section 11.7). This is because the index contents must be unaffected by default_text_search_config. If they were affected, the index contents might be inconsistent because different entries could contain tsvector that were created with different text search configurations, and there would be no way to guess which was which. It would be impossible to dump and restore such an index correctly.

Because the two-argument version of to_tsvector was used in the index above, only a query reference that uses the 2-argument version of to_tsvector with the same configuration name will use that index. That is, WHERE to_tsvector('english', body) @@ 'a & b' can use the index, but WHERE to_tsvector(body) @@ 'a & b' cannot. This ensures that an index will be used only with the same configuration used to create the index entries.

It is possible to set up more complex expression indexes wherein the configuration name is specified by another column, e.g.:

CREATE INDEX pgweb_idx ON pgweb USING GIN (to_tsvector(config_name, body));

where config_name is a column in the pgweb table. This allows mixed configurations in the same index while recording which configuration was used for each index entry. This would be useful, for example, if the document collection contained documents in different languages. Again, queries that are meant to use the index must be phrased to match, e.g., WHERE to_tsvector(config_name, body) @@ 'a & b'.

Indexes can even concatenate columns:

CREATE INDEX pgweb_idx ON pgweb USING GIN (to_tsvector('english', title || ' ' || body))
Another approach is to create a separate \texttt{tsvector} column to hold the output of \texttt{to_tsvector}. This example is a concatenation of \texttt{title} and \texttt{body}, using \texttt{coalesce} to ensure that one field will still be indexed when the other is \texttt{NULL}:

```
ALTER TABLE pgweb ADD COLUMN textsearchable_index_col tsvector;
UPDATE pgweb SET textsearchable_index_col =
    to_tsvector('english', coalesce(title,"") || ' ' || coalesce(body,""));
```

Then we create a GIN index to speed up the search:

```
CREATE INDEX textsearch_idx ON pgweb USING GIN (textsearchable_index_col);
```

Now we are ready to perform a fast full text search:

```
SELECT title
FROM pgweb
WHERE textsearchable_index_col @@ to_tsquery('create & table')
ORDER BY last_mod_date DESC
LIMIT 10;
```

When using a separate column to store the \texttt{tsvector} representation, it is necessary to create a trigger to keep the \texttt{tsvector} column current anytime \texttt{title} or \texttt{body} changes. Section 12.4.3 explains how to do that.

One advantage of the separate-column approach over an expression index is that it is not necessary to explicitly specify the text search configuration in queries in order to make use of the index. As shown in the example above, the query can depend on \texttt{default_text_search_config}. Another advantage is that searches will be faster, since it will not be necessary to redo the \texttt{to_tsvector} calls to verify index matches. (This is more important when using a GiST index than a GIN index; see Section 12.9.) The expression-index approach is simpler to set up, however, and it requires less disk space since the \texttt{tsvector} representation is not stored explicitly.

## 12.3. Controlling Text Search

To implement full text searching there must be a function to create a \texttt{tsvector} from a document and a \texttt{tsquery} from a user query. Also, we need to return results in a useful order, so we need a function that compares documents with respect to their relevance to the query. It’s also important to be able to display the results nicely. PostgreSQL provides support for all of these functions.

### 12.3.1. Parsing Documents

PostgreSQL provides the function \texttt{to_tsvector} for converting a document to the \texttt{tsvector} data type.

```
to_tsvector([ config regconfig, ] document text) returns tsvector
```

to\_tsvector parses a textual document into tokens, reduces the tokens to lexemes, and returns a \texttt{tsvector} which lists the lexemes together with their positions in the document. The document is processed according to the specified or default text search configuration. Here is a simple example:

```
SELECT to_tsvector('english’, ’a fat cat sat on a mat - it ate a fat rats’);
```
Chapter 12. Full Text Search

**to_tsvector**

> `ate`:9  `cat`:3  `fat`:2,11  `mat`:7  `rat`:12  `sat`:4

In the example above we see that the resulting **tsvector** does not contain the words `a`, `on`, or `it`, the word `rats` became `rat`, and the punctuation sign `-` was ignored.

The **to_tsvector** function internally calls a parser which breaks the document text into tokens and assigns a type to each token. For each token, a list of dictionaries (Section 12.6) is consulted, where the list can vary depending on the token type. The first dictionary that recognizes the token emits one or more normalized lexemes to represent the token. For example, `rats` became `rat` because one of the dictionaries recognized that the word `rats` is a plural form of `rat`. Some words are recognized as **stop words** (Section 12.6.1), which causes them to be ignored since they occur too frequently to be useful in searching. In our example these are `a`, `on`, and `it`. If no dictionary in the list recognizes the token then it is also ignored. In this example that happened to the punctuation sign `-` because there are in fact no dictionaries assigned for its token type (*Space symbols*), meaning space tokens will never be indexed. The choices of parser, dictionaries and which types of tokens to index are determined by the selected text search configuration (Section 12.7). It is possible to have many different configurations in the same database, and predefined configurations are available for various languages. In our example we used the default configuration **english** for the English language.

The function **setweight** can be used to label the entries of a **tsvector** with a given weight, where a weight is one of the letters **A**, **B**, **C**, or **D**. This is typically used to mark entries coming from different parts of a document, such as title versus body. Later, this information can be used for ranking of search results.

Because **to_tsvector(NULL)** will return **NULL**, it is recommended to use **coalesce** whenever a field might be null. Here is the recommended method for creating a **tsvector** from a structured document:

```sql
UPDATE tt SET ti =
    setweight(to_tsvector(coalesce(title,"")), 'A') ||
    setweight(to_tsvector(coalesce(keyword,"")), 'B') ||
    setweight(to_tsvector(coalesce(abstract,"")), 'C') ||
    setweight(to_tsvector(coalesce(body,"")), 'D');
```

Here we have used **setweight** to label the source of each lexeme in the finished **tsvector**, and then merged the labeled **tsvector** values using the **tsvector** concatenation operator **||**. (Section 12.4.1 gives details about these operations.)

### 12.3.2. Parsing Queries

PostgreSQL provides the functions **to_tquery**, **plainto_tquery**, and **phraseto_tquery** for converting a query to the **tsquery** data type. **to_tquery** offers access to more features than either **plainto_tquery** or **phraseto_tquery**, but it is less forgiving about its input.

**to_tquery** creates a **tsquery** value from **querytext**, which must consist of single tokens separated by the **tsquery** operators & (AND), | (OR), ! (NOT), and <> (FOLLOWED BY), possibly grouped using parentheses. In other words, the input to **to_tquery** must already follow the general rules for **tsquery** input, as described in Section 8.11.2. The difference is that while basic **tsquery** input takes the tokens at face value, **to_tquery** normalizes each token into a lexeme using the...
specified or default configuration, and discards any tokens that are stop words according to the configuration. For example:

```sql
SELECT to_tsquery('english', 'The & Fat & Rats');
```

```
to_tsquery
---------------
'fat' & 'rat'
```

As in basic `tsquery` input, weight(s) can be attached to each lexeme to restrict it to match only `tsvector` lexemes of those weight(s). For example:

```sql
SELECT to_tsquery('english', 'Fat | Rats:AB');
```

```
to_tsquery
------------------
'fat' | 'rat':AB
```

Also, * can be attached to a lexeme to specify prefix matching:

```sql
SELECT to_tsquery('supern:*A & star:A*B');
```

```
to_tsquery
--------------------------
'supern':*A & 'star':*AB
```

Such a lexeme will match any word in a `tsvector` that begins with the given string.

to_tsquery can also accept single-quoted phrases. This is primarily useful when the configuration includes a thesaurus dictionary that may trigger on such phrases. In the example below, a thesaurus contains the rule `supernovae stars : sn`:

```sql
SELECT to_tsquery('supernovae stars & !crab');
```

```
to_tsquery
---------------
'sn' & '!crab'
```

Without quotes, to_tsquery will generate a syntax error for tokens that are not separated by an AND, OR, or FOLLOWED BY operator.

`plainto_tsquery([ config regconfig, ] querytext text) returns tsquery` transforms the unformatted text `querytext` to a `tsquery` value. The text is parsed and normalized much as for `to_tsvector`, then the & (AND) `tsquery` operator is inserted between surviving words.

Example:

```sql
SELECT plainto_tsquery('english', 'The Fat Rats');
```

```
plainto_tsquery
-----------------
'fat' & 'rat'
```

Note that `plainto_tsquery` will not recognize `tsquery` operators, weight labels, or prefix-match labels in its input:

```sql
SELECT plainto_tsquery('english', 'The Fat & Rats:C');
```

```
plainto_tsquery
-------------------
'fat' & 'rat' & 'c'
```

Here, all the input punctuation was discarded as being space symbols.
Chapter 12. Full Text Search

phraseto_tsquery([ config regconfig, ] querytext text) returns tsquery

phraseto_tsquery behaves much like plainto_tsquery, except that it inserts the <-> (FOLLOWED BY) operator between surviving words instead of the & (AND) operator. Also, stop words are not simply discarded, but are accounted for by inserting <N> operators rather than <-> operators. This function is useful when searching for exact lexeme sequences, since the FOLLOWED BY operators check lexeme order not just the presence of all the lexemes.

Example:

```
SELECT phraseto_tsquery('english', 'The Fat Rats');
```

```
'fat' <-> 'rat'
```

Like plainto_tsquery, the phraseto_tsquery function will not recognize tsquery operators, weight labels, or prefix-match labels in its input:

```
SELECT phraseto_tsquery('english', 'The Fat & Rats:C');
```

```
'fat' <-> 'rat' <-> 'c'
```

12.3.3. Ranking Search Results

Ranking attempts to measure how relevant documents are to a particular query, so that when there are many matches the most relevant ones can be shown first. PostgreSQL provides two predefined ranking functions, which take into account lexical, proximity, and structural information; that is, they consider how often the query terms appear in the document, how close together the terms are in the document, and how important is the part of the document where they occur. However, the concept of relevancy is vague and very application-specific. Different applications might require additional information for ranking, e.g., document modification time. The built-in ranking functions are only examples. You can write your own ranking functions and/or combine their results with additional factors to fit your specific needs.

The two ranking functions currently available are:

```
ts_rank([ weights float4[], ] vector tsvector, query tsquery [ , normalization integer ] ) returns float4
```

Ranks vectors based on the frequency of their matching lexemes.

```
ts_rank_cd([ weights float4[], ] vector tsvector, query tsquery [ , normalization integer ] ) returns float4
```

This function computes the cover density ranking for the given document vector and query, as described in Clarke, Cormack, and Tudhope’s "Relevance Ranking for One to Three Term Queries" in the journal "Information Processing and Management", 1999. Cover density is similar to ts_rank ranking except that the proximity of matching lexemes to each other is taken into consideration.

This function requires lexeme positional information to perform its calculation. Therefore, it ignores any “stripped” lexemes in the tsvector. If there are no unstripped lexemes in the input,
Chapter 12. Full Text Search

the result will be zero. (See Section 12.4.1 for more information about the strip function and positional information in tsvectors.)

For both these functions, the optional weights argument offers the ability to weigh word instances more or less heavily depending on how they are labeled. The weight arrays specify how heavily to weigh each category of word, in the order:

{D-weight, C-weight, B-weight, A-weight}

If no weights are provided, then these defaults are used:

{0.1, 0.2, 0.4, 1.0}

Typically weights are used to mark words from special areas of the document, like the title or an initial abstract, so they can be treated with more or less importance than words in the document body.

Since a longer document has a greater chance of containing a query term it is reasonable to take into account document size, e.g., a hundred-word document with five instances of a search word is probably more relevant than a thousand-word document with five instances. Both ranking functions take an integer normalization option that specifies whether and how a document’s length should impact its rank. The integer option controls several behaviors, so it is a bit mask: you can specify one or more behaviors using | (for example, 2|4).

- 0 (the default) ignores the document length
- 1 divides the rank by 1 + the logarithm of the document length
- 2 divides the rank by the document length
- 4 divides the rank by the mean harmonic distance between extents (this is implemented only by ts_rank_cd)
- 8 divides the rank by the number of unique words in document
- 16 divides the rank by 1 + the logarithm of the number of unique words in document
- 32 divides the rank by itself + 1

If more than one flag bit is specified, the transformations are applied in the order listed.

It is important to note that the ranking functions do not use any global information, so it is impossible to produce a fair normalization to 1% or 100% as sometimes desired. Normalization option 32 (rank/(rank+1)) can be applied to scale all ranks into the range zero to one, but of course this is just a cosmetic change; it will not affect the ordering of the search results.

Here is an example that selects only the ten highest-ranked matches:

```
SELECT title, ts_rank_cd(textsearch, query) AS rank
FROM apod, to_tsquery('neutrino|(dark & matter)') query
WHERE query @@ textsearch
ORDER BY rank DESC
LIMIT 10;
```

<table>
<thead>
<tr>
<th>title</th>
<th>rank</th>
</tr>
</thead>
<tbody>
<tr>
<td>Neutrinos in the Sun</td>
<td>3.1</td>
</tr>
<tr>
<td>The Sudbury Neutrino Detector</td>
<td>2.4</td>
</tr>
<tr>
<td>A MACHO View of Galactic Dark Matter</td>
<td>2.01317</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>1.91171</td>
</tr>
<tr>
<td>The Virgo Cluster: Hot Plasma and Dark Matter</td>
<td>1.90953</td>
</tr>
<tr>
<td>Rafting for Solar Neutrinos</td>
<td>1.9</td>
</tr>
<tr>
<td>NGC 4650A: Strange Galaxy and Dark Matter</td>
<td>1.85774</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>1.6123</td>
</tr>
</tbody>
</table>
Chapter 12. Full Text Search

This is the same example using normalized ranking:

```sql
SELECT title, ts_rank_cd(textsearch, query, 32 /* rank/(rank+1) */ ) AS rank
FROM apod, to_tsquery('neutrino|(dark & matter)') query
WHERE query @@ textsearch
ORDER BY rank DESC
LIMIT 10;
```

<table>
<thead>
<tr>
<th>title</th>
<th>rank</th>
</tr>
</thead>
<tbody>
<tr>
<td>Neutrinos in the Sun</td>
<td>0.756097569485493</td>
</tr>
<tr>
<td>The Sudbury Neutrino Detector</td>
<td>0.705882361190954</td>
</tr>
<tr>
<td>A MACHO View of Galactic Dark Matter</td>
<td>0.668123210574724</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>0.65655958650282</td>
</tr>
<tr>
<td>The Virgo Cluster: Hot Plasma and Dark Matter</td>
<td>0.656301290640973</td>
</tr>
<tr>
<td>Rafting for Solar Neutrinos</td>
<td>0.655172410958162</td>
</tr>
<tr>
<td>NGC 4650A: Strange Galaxy and Dark Matter</td>
<td>0.650072921219637</td>
</tr>
<tr>
<td>Hot Gas and Dark Matter</td>
<td>0.617195790024749</td>
</tr>
<tr>
<td>Ice Fishing for Cosmic Neutrinos</td>
<td>0.615384618911517</td>
</tr>
<tr>
<td>Weak Lensing Distorts the Universe</td>
<td>0.450010798361481</td>
</tr>
</tbody>
</table>

Ranking can be expensive since it requires consulting the tsvector of each matching document, which can be I/O bound and therefore slow. Unfortunately, it is almost impossible to avoid since practical queries often result in large numbers of matches.

### 12.3.4. Highlighting Results

To present search results it is ideal to show a part of each document and how it is related to the query. Usually, search engines show fragments of the document with marked search terms. PostgreSQL provides a function `ts_headline` that implements this functionality.

`ts_headline([ config regconfig, ] document text, query tsquery [, options text ])` returns text

`ts_headline` accepts a document along with a query, and returns an excerpt from the document in which terms from the query are highlighted. The configuration to be used to parse the document can be specified by `config`; if `config` is omitted, the `default_text_search_config` configuration is used.

If an `options` string is specified it must consist of a comma-separated list of one or more `option=value` pairs. The available options are:

- `StartSel, StopSel`: the strings with which to delimit query words appearing in the document, to distinguish them from other excerpted words. You must double-quote these strings if they contain spaces or commas.
- `MaxWords, MinWords`: these numbers determine the longest and shortest headlines to output.
- `ShortWord`: words of this length or less will be dropped at the start and end of a headline. The default value of three eliminates common English articles.
- `HighlightAll`: Boolean flag; if `true` the whole document will be used as the headline, ignoring the preceding three parameters.
- `MaxFragments`: maximum number of text excerpts or fragments to display. The default value of zero selects a non-fragment-oriented headline generation method. A value greater than zero selects
fragment-based headline generation. This method finds text fragments with as many query words as possible and stretches those fragments around the query words. As a result query words are close to the middle of each fragment and have words on each side. Each fragment will be of at most \text{MaxWords} and words of length \text{ShortWord} or less are dropped at the start and end of each fragment. If not all query words are found in the document, then a single fragment of the first \text{MinWords} in the document will be displayed.

- **FragmentDelimiter**: When more than one fragment is displayed, the fragments will be separated by this string.

Any unspecified options receive these defaults:

\begin{verbatim}
StartSel=<b>, StopSel=<b>,
MaxWords=35, MinWords=15, ShortWord=3, HighlightAll=FALSE,
MaxFragments=0, FragmentDelimiter=" ... "
\end{verbatim}

For example:

\begin{verbatim}
SELECT ts_headline('english',
  'The most common type of search is to find all documents containing given query terms and return them in order of their similarity to the query.',
  to_tsquery('query & similarity'));
\end{verbatim}

\texttt{ts_headline}

\begin{verbatim}
(containing given \texttt{query} terms and return them in order of their \texttt{similarity} to the \texttt{query}).
\end{verbatim}

\begin{verbatim}
SELECT ts_headline('english',
  'The most common type of search is to find all documents containing given query terms and return them in order of their similarity to the query.',
  to_tsquery('query & similarity'),
  'StartSel = <, StopSel = >');
\end{verbatim}

\texttt{ts_headline}

\begin{verbatim}
(containing given \texttt{query} terms and return them in order of their \texttt{similarity} to the \texttt{query}).
\end{verbatim}

\texttt{ts_headline} uses the original document, not a \texttt{tsvector} summary, so it can be slow and should be used with care.

### 12.4. Additional Features

This section describes additional functions and operators that are useful in connection with text search.
12.4.1. Manipulating Documents

Section 12.3.1 showed how raw textual documents can be converted into \textit{tsvector} values. PostgreSQL also provides functions and operators that can be used to manipulate documents that are already in \textit{tsvector} form.

\texttt{tsvector \|\| tsvector}

The \texttt{tsvector} concatenation operator returns a vector which combines the lexemes and positional information of the two vectors given as arguments. Positions and weight labels are retained during the concatenation. Positions appearing in the right-hand vector are offset by the largest position mentioned in the left-hand vector, so that the result is nearly equivalent to the result of performing \texttt{to_tsvector} on the concatenation of the two original document strings. (The equivalence is not exact, because any stop-words removed from the end of the left-hand argument will not affect the result, whereas they would have affected the positions of the lexemes in the right-hand argument if textual concatenation were used.)

One advantage of using concatenation in the vector form, rather than concatenating text before applying \texttt{to_tsvector}, is that you can use different configurations to parse different sections of the document. Also, because the \texttt{setweight} function marks all lexemes of the given vector the same way, it is necessary to parse the text and do \texttt{setweight} before concatenating if you want to label different parts of the document with different weights.

\texttt{setweight(vector tsvector, weight \texttt{"char"}) returns tsvector}

\texttt{setweight} returns a copy of the input vector in which every position has been labeled with the given \texttt{weight}, either A, B, C, or D. (D is the default for new vectors and as such is not displayed on output.) These labels are retained when vectors are concatenated, allowing words from different parts of a document to be weighted differently by ranking functions.

Note that weight labels apply to \textit{positions}, not \textit{lexemes}. If the input vector has been stripped of positions then \texttt{setweight} does nothing.

\texttt{length(vector tsvector) returns integer}

Returns the number of lexemes stored in the vector.

\texttt{strip(vector tsvector) returns tsvector}

Returns a vector that lists the same lexemes as the given vector, but lacks any position or weight information. The result is usually much smaller than an unstripped vector, but it is also less useful. Relevance ranking does not work as well on stripped vectors as unstripped ones. Also, the \texttt{\textless\textgreater} (FOLLOWED BY) \texttt{tsquery} operator will never match stripped input, since it cannot determine the distance between lexeme occurrences.

A full list of \textit{tsvector}-related functions is available in Table 9-40.

12.4.2. Manipulating Queries

Section 12.3.2 showed how raw textual queries can be converted into \textit{tsquery} values. PostgreSQL also provides functions and operators that can be used to manipulate queries that are already in \textit{tsquery} form.

\texttt{tsquery \&\& tsquery}

Returns the AND-combination of the two given queries.
tsquery || tsquery
   Returns the OR-combination of the two given queries.

!! tsquery
   Returns the negation (NOT) of the given query.

tquery <--> tsquery
   Returns a query that searches for a match to the first given query immediately followed by a
   match to the second given query, using the <--> (FOLLOWED BY) tsquery operator. For example:
   SELECT to_tsquery('fat') <--> to_tsquery('cat | rat');
   -----------------------------------------------
   'fat' <--> 'cat' | 'fat' <--> 'rat'

tquery_phrase(query1 tsquery, query2 tsquery [, distance integer ]) returns tsquery
   Returns a query that searches for a match to the first given query followed by a match to the
   second given query at a distance of at distance lexemes, using the <N> tsquery operator.
   For example:
   SELECT tsquery_phrase(to_tsquery('fat'), to_tsquery('cat'), 10);
   tsquery_phrase
   ------------------
   'fat' <10> 'cat'

numnode(query tsquery) returns integer
   Returns the number of nodes (lexemes plus operators) in a tsquery. This function is useful to
determine if the query is meaningful (returns > 0), or contains only stop words (returns 0).
   Examples:
   SELECT numnode(plainto_tsquery('the any'));
   NOTICE: query contains only stopword(s) or doesn’t contain lexeme(s), ignored
   numnode
   ---------
   0
   SELECT numnode('foo & bar'::tsquery);
   numnode
   ---------
   3

querytree(query tsquery) returns text
   Returns the portion of a tsquery that can be used for searching an index. This function is useful
   for detecting unindexable queries, for example those containing only stop words or only negated
terms. For example:
   SELECT querytree(to_tsquery('!defined'));
   querytree
   12.4.2.1. Query Rewriting

The ts_rewrite family of functions search a given tsquery for occurrences of a target subquery,
and replace each occurrence with a substitute subquery. In essence this operation is a tsquery-
specific version of substring replacement. A target and substitute combination can be thought of as a
query rewrite rule. A collection of such rewrite rules can be a powerful search aid. For example, you can expand the search using synonyms (e.g., new york, big apple, nyc, gotham) or narrow the search to direct the user to some hot topic. There is some overlap in functionality between this feature and thesaurus dictionaries (Section 12.6.4). However, you can modify a set of rewrite rules on-the-fly without reindexing, whereas updating a thesaurus requires reindexing to be effective.

\[
\text{ts\_rewrite (query tsquery, target tsquery, substitute tsquery) returns tsquery}
\]

This form of \text{ts\_rewrite} simply applies a single rewrite rule: target is replaced by substitute wherever it appears in query. For example:

\[
\text{SELECT ts\_rewrite('a & b'::tsquery, 'a'::tsquery, 'c'::tsquery)};
\]

\[
\text{ts\_rewrite}
\]

\[
'b' & 'c'
\]

\[
\text{ts\_rewrite (query tsquery, select text) returns tsquery}
\]

This form of \text{ts\_rewrite} accepts a starting query and a SQL select command, which is given as a text string. The select must yield two columns of tsquery type. For each row of the select result, occurrences of the first column value (the target) are replaced by the second column value (the substitute) within the current query value. For example:

\[
\text{CREATE TABLE aliases (t tsquery primary key, s tsquery)};
\]

\[
\text{INSERT INTO aliases VALUES('a', 'c')};
\]

\[
\text{SELECT ts\_rewrite('a & b'::tsquery, 'SELECT t,s FROM aliases')};
\]

\[
\text{ts\_rewrite}
\]

\[
'b' & 'c'
\]

Note that when multiple rewrite rules are applied in this way, the order of application can be important; so in practice you will want the source query to \text{ORDER BY} some ordering key.

Let’s consider a real-life astronomical example. We’ll expand query supernovae using table-driven rewriting rules:

\[
\text{CREATE TABLE aliases (t tsquery primary key, s tsquery)};
\]

\[
\text{INSERT INTO aliases VALUES(to\_tsquery('supernovae'), to\_tsquery('supernovae|sn'))};
\]

\[
\text{SELECT ts\_rewrite(to\_tsquery('supernovae & crab'), 'SELECT * FROM aliases')};
\]

\[
\text{ts\_rewrite}
\]

\[
'crab' & ( 'supernova' | 'sn' )
\]

We can change the rewriting rules just by updating the table:

\[
\text{UPDATE aliases}
\]

\[
\text{SET s = to\_tsquery('supernovae|sn & !nebulae')};
\]

\[
\text{WHERE t = to\_tsquery('supernovae')};
\]

\[
\text{SELECT ts\_rewrite(to\_tsquery('supernovae & crab'), 'SELECT * FROM aliases')};
\]

\[
\text{ts\_rewrite}
\]

\[
'crab' & ( 'supernova' | 'sn' & !'nebula' )
\]

Rewriting can be slow when there are many rewriting rules, since it checks every rule for a possible match. To filter out obvious non-candidate rules we can use the containment operators for the
tsquery type. In the example below, we select only those rules which might match the original query:

```sql
SELECT ts_rewrite('a & b'::tsquery,
    'SELECT t,s FROM aliases WHERE "a & b"::tsquery @> t');
```

```
ts_rewrite
----------------
'b' & 'c'
```

12.4.3. Triggers for Automatic Updates

When using a separate column to store the tsvector representation of your documents, it is necessary to create a trigger to update the tsvector column when the document content columns change. Two built-in trigger functions are available for this, or you can write your own.

```sql
tsvector_update_trigger(tsvector_column_name, config_name, text_column_name [, ... ])
tsvector_update_trigger_column(tsvector_column_name, config_column_name, text_column_name [, ... ])
```

These trigger functions automatically compute a tsvector column from one or more textual columns, under the control of parameters specified in the CREATE TRIGGER command. An example of their use is:

```sql
CREATE TABLE messages (    title text,    body text,    tsv tsvector );
CREATE TRIGGER tsvectorupdate BEFORE INSERT OR UPDATE ON messages FOR EACH ROW EXECUTE PROCEDURE tsvector_update_trigger(tsv, 'pg_catalog.english', title, body);
```

```sql
INSERT INTO messages VALUES('title here', 'the body text is here');
```

```sql
SELECT * FROM messages;
```

```
<table>
<thead>
<tr>
<th>title</th>
<th>body</th>
<th>tsv</th>
</tr>
</thead>
<tbody>
<tr>
<td>title here</td>
<td>the body text is here</td>
<td>'bodi':4 'text':5 'titl':1</td>
</tr>
</tbody>
</table>
```

```sql
SELECT title, body FROM messages WHERE tsv @@ to_tsquery('title & body');
```

```
<table>
<thead>
<tr>
<th>title</th>
<th>body</th>
</tr>
</thead>
<tbody>
<tr>
<td>title here</td>
<td>the body text is here</td>
</tr>
</tbody>
</table>
```

Having created this trigger, any change in title or body will automatically be reflected into tsv, without the application having to worry about it.

The first trigger argument must be the name of the tsvector column to be updated. The second argument specifies the text search configuration to be used to perform the conversion. For tsvector_update_trigger, the configuration name is simply given as the second trigger argument. It must be schema-qualified as shown above, so that the trigger behavior will not change with changes in search_path. For tsvector_update_trigger_column, the second trigger argument is the name of another table column, which must be of type regconfig. This allows a
per-row selection of configuration to be made. The remaining argument(s) are the names of textual columns (of type text, varchar, or char). These will be included in the document in the order given. NULL values will be skipped (but the other columns will still be indexed).

A limitation of these built-in triggers is that they treat all the input columns alike. To process columns differently — for example, to weight title differently from body — it is necessary to write a custom trigger. Here is an example using PL/pgSQL as the trigger language:

```sql
CREATE FUNCTION messages_trigger() RETURNS trigger AS $$
begin
    new.tsv :=
        setweight(to_tsvector('pg_catalog.english', coalesce(new.title,"")), 'A') ||
        setweight(to_tsvector('pg_catalog.english', coalesce(new.body,"")), 'D');
    return new;
end $$ LANGUAGE plpgsql;

CREATE TRIGGER tsvectorupdate BEFORE INSERT OR UPDATE
    ON messages FOR EACH ROW EXECUTE PROCEDURE messages_trigger();
```

Keep in mind that it is important to specify the configuration name explicitly when creating tsvector values inside triggers, so that the column’s contents will not be affected by changes to default_text_search_config. Failure to do this is likely to lead to problems such as search results changing after a dump and reload.

### 12.4.4. Gathering Document Statistics

The function `ts_stat` is useful for checking your configuration and for finding stop-word candidates.

```sql
ts_stat(sqlquery text, [ weights text, ]
    OUT word text, OUT ndoc integer,
    OUT nentry integer) returns setof record
```

`sqlquery` is a text value containing an SQL query which must return a single tsvector column. `ts_stat` executes the query and returns statistics about each distinct lexeme (word) contained in the tsvector data. The columns returned are

- `word text` — the value of a lexeme
- `ndoc integer` — number of documents (tsvectors) the word occurred in
- `nentry integer` — total number of occurrences of the word

If `weights` is supplied, only occurrences having one of those weights are counted.

For example, to find the ten most frequent words in a document collection:

```sql
SELECT * FROM ts_stat('SELECT vector FROM apod')
ORDER BY nentry DESC, ndoc DESC, word
LIMIT 10;
```

The same, but counting only word occurrences with weight A or B:

```sql
SELECT * FROM ts_stat('SELECT vector FROM apod', 'ab')
ORDER BY nentry DESC, ndoc DESC, word
```
12.5. Parsers

Text search parsers are responsible for splitting raw document text into *tokens* and identifying each token's type, where the set of possible types is defined by the parser itself. Note that a parser does not modify the text at all — it simply identifies plausible word boundaries. Because of this limited scope, there is less need for application-specific custom parsers than there is for custom dictionaries. At present PostgreSQL provides just one built-in parser, which has been found to be useful for a wide range of applications.

The built-in parser is named `pg_catalog.default`. It recognizes 23 token types, shown in Table 12-1.

Table 12-1. Default Parser’s Token Types

<table>
<thead>
<tr>
<th>Alias</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>Word, all ASCII letters</td>
<td>elephant</td>
</tr>
<tr>
<td>word</td>
<td>Word, all letters</td>
<td>mañana</td>
</tr>
<tr>
<td>numword</td>
<td>Word, letters and digits</td>
<td>beta1</td>
</tr>
<tr>
<td>asciihword</td>
<td>Hyphenated word, all ASCII</td>
<td>up-to-date</td>
</tr>
<tr>
<td>hword</td>
<td>Hyphenated word, all letters</td>
<td>lógico-matemática</td>
</tr>
<tr>
<td>numhword</td>
<td>Hyphenated word, letters and digits</td>
<td>postgresql-beta1</td>
</tr>
<tr>
<td>hword_asciipart</td>
<td>Hyphenated word part, all ASCII</td>
<td>postgresql in the context</td>
</tr>
<tr>
<td>hword_part</td>
<td>Hyphenated word part, all letters</td>
<td>lógico or matemática in the context</td>
</tr>
<tr>
<td>hword_numpart</td>
<td>Hyphenated word part, letters and digits</td>
<td>beta1 in the context</td>
</tr>
<tr>
<td>email</td>
<td>Email address</td>
<td><a href="mailto:foo@example.com">foo@example.com</a></td>
</tr>
<tr>
<td>protocol</td>
<td>Protocol head</td>
<td>http://</td>
</tr>
<tr>
<td>url</td>
<td>URL</td>
<td>example.com/stuff/index.html</td>
</tr>
<tr>
<td>host</td>
<td>Host</td>
<td>example.com</td>
</tr>
<tr>
<td>url_path</td>
<td>URL path</td>
<td>/stuff/index.html, in the context of a URL</td>
</tr>
<tr>
<td>file</td>
<td>File or path name</td>
<td>/usr/local/foo.txt, if not within a URL</td>
</tr>
<tr>
<td>sfloat</td>
<td>Scientific notation</td>
<td>-1.234e56</td>
</tr>
<tr>
<td>float</td>
<td>Decimal notation</td>
<td>-1.234</td>
</tr>
<tr>
<td>int</td>
<td>Signed integer</td>
<td>-1234</td>
</tr>
<tr>
<td>uint</td>
<td>Unsigned integer</td>
<td>1234</td>
</tr>
<tr>
<td>version</td>
<td>Version number</td>
<td>8.3.0</td>
</tr>
</tbody>
</table>
### 12.6. Dictionaries

Dictionaries are used to eliminate words that should not be considered in a search (*stop words*), and to *normalize* words so that different derived forms of the same word will match. A successfully normalized word is called a *lexeme*. Aside from improving search quality, normalization and removal of stop words reduce the size of the *tsvector* representation of a document, thereby improving performance. Normalization does not always have linguistic meaning and usually depends on application semantics.

Some examples of normalization:

<table>
<thead>
<tr>
<th>Alias</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>tag</td>
<td>XML tag</td>
<td><code>&lt;a href=&quot;dictionaries.html&quot;&gt;</code></td>
</tr>
<tr>
<td>entity</td>
<td>XML entity</td>
<td><code>&amp;amp;</code></td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td>(any whitespace or punctuation not otherwise recognized)</td>
</tr>
</tbody>
</table>

Note: The parser’s notion of a “letter” is determined by the database’s locale setting, specifically `lc_type`. Words containing only the basic ASCII letters are reported as a separate token type, since it is sometimes useful to distinguish them. In most European languages, token types `word` and `asciword` should be treated alike.

(email) does not support all valid email characters as defined by RFC 5322. Specifically, the only non-alphanumeric characters supported for email user names are period, dash, and underscore.

It is possible for the parser to produce overlapping tokens from the same piece of text. As an example, a hyphenated word will be reported both as the entire word and as each component:

```
SELECT alias, description, token FROM ts_debug('foo-bar-beta1');
```

<table>
<thead>
<tr>
<th>alias</th>
<th>description</th>
<th>token</th>
</tr>
</thead>
<tbody>
<tr>
<td>numhword</td>
<td>Hyphenated word, letters and digits</td>
<td>foo-bar-beta1</td>
</tr>
<tr>
<td>hword_asciipart</td>
<td>Hyphenated word part, all ASCII</td>
<td>foo</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td>-</td>
</tr>
<tr>
<td>hword_asciipart</td>
<td>Hyphenated word part, all ASCII</td>
<td>bar</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td>-</td>
</tr>
<tr>
<td>hword_numpart</td>
<td>Hyphenated word part, letters and digits</td>
<td>betal</td>
</tr>
</tbody>
</table>

This behavior is desirable since it allows searches to work for both the whole compound word and for components. Here is another instructive example:

```
SELECT alias, description, token FROM ts_debug('http://example.com/stuff/index.html');
```

<table>
<thead>
<tr>
<th>alias</th>
<th>description</th>
<th>token</th>
</tr>
</thead>
<tbody>
<tr>
<td>protocol</td>
<td>Protocol head</td>
<td>http://</td>
</tr>
<tr>
<td>url</td>
<td>URL</td>
<td>example.com/stuff/index.html</td>
</tr>
<tr>
<td>host</td>
<td>Host</td>
<td>example.com</td>
</tr>
<tr>
<td>url_path</td>
<td>URL path</td>
<td>/stuff/index.html</td>
</tr>
</tbody>
</table>

12.6. Dictionaries
• Linguistic - Ispell dictionaries try to reduce input words to a normalized form; stemmer dictionaries remove word endings
• URL locations can be canonicalized to make equivalent URLs match:
  • http://www.pgsql.ru/db/mw/index.html
  • http://www.pgsql.ru/db/mw/
  • http://www.pgsql.ru/db/../db/mw/index.html
• Color names can be replaced by their hexadecimal values, e.g., red, green, blue, magenta
  -> FF0000, 00FF00, 0000FF, FF00FF
• If indexing numbers, we can remove some fractional digits to reduce the range of possible numbers, so for example 3.14159265359, 3.1415926, 3.14 will be the same after normalization if only two digits are kept after the decimal point.

A dictionary is a program that accepts a token as input and returns:

• an array of lexemes if the input token is known to the dictionary (notice that one token can produce more than one lexeme)
• a single lexeme with the TSL_FILTER flag set, to replace the original token with a new token to be passed to subsequent dictionaries (a dictionary that does this is called a filtering dictionary)
• an empty array if the dictionary knows the token, but it is a stop word
• NULL if the dictionary does not recognize the input token

PostgreSQL provides predefined dictionaries for many languages. There are also several predefined templates that can be used to create new dictionaries with custom parameters. Each predefined dictionary template is described below. If no existing template is suitable, it is possible to create new ones; see the contrib/ area of the PostgreSQL distribution for examples.

A text search configuration binds a parser together with a set of dictionaries to process the parser’s output tokens. For each token type that the parser can return, a separate list of dictionaries is specified by the configuration. When a token of that type is found by the parser, each dictionary in the list is consulted in turn, until some dictionary recognizes it as a known word. If it is identified as a stop word, or if no dictionary recognizes the token, it will be discarded and not indexed or searched for. Normally, the first dictionary that returns a non-NULL output determines the result, and any remaining dictionaries are not consulted; but a filtering dictionary can replace the given word with a modified word, which is then passed to subsequent dictionaries.

The general rule for configuring a list of dictionaries is to place first the most narrow, most specific dictionary, then the more general dictionaries, finishing with a very general dictionary, like a Snowball stemmer or simple, which recognizes everything. For example, for an astronomy-specific search (astro_en configuration) one could bind token type asciiword (ASCII word) to a synonym dictionary of astronomical terms, a general English dictionary and a Snowball English stemmer:

```sql
ALTER TEXT SEARCH CONFIGURATION astro_en
    ADD MAPPING FOR asciiword WITH astrosyn, english_ispell, english_stem;
```

A filtering dictionary can be placed anywhere in the list, except at the end where it’d be useless. Filtering dictionaries are useful to partially normalize words to simplify the task of later dictionaries. For example, a filtering dictionary could be used to remove accents from accented letters, as is done by the unaccent module.
12.6.1. Stop Words

Stop words are words that are very common, appear in almost every document, and have no discrimination value. Therefore, they can be ignored in the context of full text searching. For example, every English text contains words like a and the, so it is useless to store them in an index. However, stop words do affect the positions in tsvector, which in turn affect ranking:

```sql
SELECT to_tsvector('english','in the list of stop words');
```

```
           ---------------------------
           |  'list':3 'stop':5 'word':6 |
```

The missing positions 1,2,4 are because of stop words. Ranks calculated for documents with and without stop words are quite different:

```sql
SELECT ts_rank_cd (to_tsvector('english','in the list of stop words'), to_tsquery('list & stop'));
```

```
   ts_rank_cd
------------
       0.05
```

```sql
SELECT ts_rank_cd (to_tsvector('english','list stop words'), to_tsquery('list & stop'));
```

```
   ts_rank_cd
------------
       0.1
```

It is up to the specific dictionary how it treats stop words. For example, ispell dictionaries first normalize words and then look at the list of stop words, while Snowball stemmers first check the list of stop words. The reason for the different behavior is an attempt to decrease noise.

12.6.2. Simple Dictionary

The simple dictionary template operates by converting the input token to lower case and checking it against a file of stop words. If it is found in the file then an empty array is returned, causing the token to be discarded. If not, the lower-cased form of the word is returned as the normalized lexeme. Alternatively, the dictionary can be configured to report non-stop-words as unrecognized, allowing them to be passed on to the next dictionary in the list.

Here is an example of a dictionary definition using the simple template:

```sql
CREATE TEXT SEARCH DICTIONARY public.simple_dict (
    TEMPLATE = pg_catalog.simple,
    STOPWORDS = english
);
```

Here, english is the base name of a file of stop words. The file’s full name will be $SHAREDIR/tsearch_data/english.stop, where $SHAREDIR means the PostgreSQL installation’s shared-data directory, often /usr/local/share/postgresql (use pg_config --sharedir to determine it if you’re not sure). The file format is simply a list of words, one per line. Blank lines and trailing spaces are ignored, and upper case is folded to lower case, but no other processing is done on the file contents.

Now we can test our dictionary:

```sql
SELECT ts_lexize('public.simple_dict','YeS');
```
Chapter 12. Full Text Search

```
ts_lexize
-----------
{yes}

SELECT ts_lexize('public.simple_dict','The');
ts_lexize
-----------
{}

We can also choose to return NULL, instead of the lower-cased word, if it is not found in the stop words file. This behavior is selected by setting the dictionary’s Accept parameter to false. Continuing the example:

ALTER TEXT SEARCH DICTIONARY public.simple_dict ( Accept = false );

SELECT ts_lexize('public.simple_dict','YeS');
ts_lexize
-----------
{}

SELECT ts_lexize('public.simple_dict','The');
ts_lexize
-----------
{}

With the default setting of Accept = true, it is only useful to place a simple dictionary at the end of a list of dictionaries, since it will never pass on any token to a following dictionary. Conversely, Accept = false is only useful when there is at least one following dictionary.

Caution
---
Most types of dictionaries rely on configuration files, such as files of stop words. These files must be stored in UTF-8 encoding. They will be translated to the actual database encoding, if that is different, when they are read into the server.

Caution
---
Normally, a database session will read a dictionary configuration file only once, when it is first used within the session. If you modify a configuration file and want to force existing sessions to pick up the new contents, issue an ALTER TEXT SEARCH DICTIONARY command on the dictionary. This can be a “dummy” update that doesn’t actually change any parameter values.

12.6.3. Synonym Dictionary

This dictionary template is used to create dictionaries that replace a word with a synonym. Phrases are not supported (use the thesaurus template (Section 12.6.4) for that). A synonym dictionary can be used to overcome linguistic problems, for example, to prevent an English stemmer dictionary
from reducing the word “Paris” to “pari”. It is enough to have a Paris pari line in the synonym dictionary and put it before the english_stem dictionary. For example:

```
SELECT * FROM ts_debug('english', 'Paris');
alias | description | token | dictionaries | dictionary | lexemes
-----------------+-----------------+-------+----------------+--------------+---------
asciiword | Word, all ASCII | Paris | {english_stem} | english_stem | {pari}
```

CREATE TEXT SEARCH DICTIONARY my_synonym {
    TEMPLATE = synonym,
    SYNONYMS = my_synonyms
};

ALTER TEXT SEARCH CONFIGURATION english
    ALTER MAPPING FOR asciiword
    WITH my_synonym, english_stem;

SELECT * FROM ts_debug('english', 'Paris');
alias | description | token | dictionaries | dictionary | lexemes
-----------------+-----------------+-------+---------------------------+------------+---------
asciiword | Word, all ASCII | Paris | {my_synonym,english_stem} | my_synonym | {paris}
```

The only parameter required by the synonym template is SYNONYMS, which is the base name of its configuration file — my_synonyms in the above example. The file’s full name will be $SHAREDIR/tsearch_data/my_synonyms.syn (where $SHAREDIR means the PostgreSQL installation’s shared-data directory). The file format is just one line per word to be substituted, with the word followed by its synonym, separated by white space. Blank lines and trailing spaces are ignored.

The synonym template also has an optional parameter CaseSensitive, which defaults to false. When CaseSensitive is false, words in the synonym file are folded to lower case, as are input tokens. When it is true, words and tokens are not folded to lower case, but are compared as-is.

An asterisk (*) can be placed at the end of a synonym in the configuration file. This indicates that the synonym is a prefix. The asterisk is ignored when the entry is used in to_tsvector(), but when it is used in to_tsquery(), the result will be a query item with the prefix match marker (see Section 12.3.2). For example, suppose we have these entries in $SHAREDIR/tsearch_data/synonym_sample.syn:

```
postgres pgsql
postgresql pgsql
postgre pgsql
google googl
indices index*
```

Then we will get these results:

```
mydb=# CREATE TEXT SEARCH DICTIONARY syn (template=synonym, synonyms='synonym_sample');
mydb=# SELECT ts_lexize('syn','indices');
ts_lexize
-----------------+-------+---------
{index}          (1 row)
```

mydb=# CREATE TEXT SEARCH CONFIGURATION tst (copy=simple);
mydb=# ALTER TEXT SEARCH CONFIGURATION tst ALTER MAPPING FOR asciiword WITH syn;

398
mydb=# SELECT to_tsvector('tst','indices');
  to_tsvector
-------------
  'index':1
(1 row)

mydb=# SELECT to_tsquery('tst','indices');
  to_tsquery
-----------
  'index':*
(1 row)

mydb=# SELECT 'indexes are very useful'::tsvector;
  tsvector
---------------------------------
  'are' 'indexes' 'useful' 'very'
(1 row)

mydb=# SELECT 'indexes are very useful'::tsvector @@ to_tsquery('tst','indices');
?column?
--------
  t
(1 row)

12.6.4. Thesaurus Dictionary

A thesaurus dictionary (sometimes abbreviated as TZ) is a collection of words that includes information about the relationships of words and phrases, i.e., broader terms (BT), narrower terms (NT), preferred terms, non-preferred terms, related terms, etc.

Basically a thesaurus dictionary replaces all non-preferred terms by one preferred term and, optionally, preserves the original terms for indexing as well. PostgreSQL’s current implementation of the thesaurus dictionary is an extension of the synonym dictionary with added phrase support. A thesaurus dictionary requires a configuration file of the following format:

```
# this is a comment
sample word(s) : indexed word(s)
more sample word(s) : more indexed word(s)
...
```

where the colon (:) symbol acts as a delimiter between a phrase and its replacement.

A thesaurus dictionary uses a subdictionary (which is specified in the dictionary’s configuration) to normalize the input text before checking for phrase matches. It is only possible to select one subdictionary. An error is reported if the subdictionary fails to recognize a word. In that case, you should remove the use of the word or teach the subdictionary about it. You can place an asterisk (*) at the beginning of an indexed word to skip applying the subdictionary to it, but all sample words must be known to the subdictionary.

The thesaurus dictionary chooses the longest match if there are multiple phrases matching the input, and ties are broken by using the last definition.
Specific stop words recognized by the subdictionary cannot be specified; instead use \( ? \) to mark the location where any stop word can appear. For example, assuming that \( a \) and \( the \) are stop words according to the subdictionary:

\[
? \text{one} \ ? \text{two} : \text{swsw}
\]

matches \( a \text{one the two} \) and \( \text{the one a two} \); both would be replaced by \( \text{swsw} \).

Since a thesaurus dictionary has the capability to recognize phrases it must remember its state and interact with the parser. A thesaurus dictionary uses these assignments to check if it should handle the next word or stop accumulation. The thesaurus dictionary must be configured carefully. For example, if the thesaurus dictionary is assigned to handle only the \( \text{asciiword} \) token, then a thesaurus dictionary definition like \( \text{one 7} \) will not work since token type \( \text{uint} \) is not assigned to the thesaurus dictionary.

### Caution

Thesauruses are used during indexing so any change in the thesaurus dictionary's parameters requires reindexing. For most other dictionary types, small changes such as adding or removing stopwords does not force reindexing.

#### 12.6.4.1. Thesaurus Configuration

To define a new thesaurus dictionary, use the \( \text{thesaurus} \) template. For example:

```
CREATE TEXT SEARCH DICTIONARY thesaurus_simple (
    TEMPLATE = thesaurus,
    DictFile = mythesaurus,
    Dictionary = pg_catalog.english_stem
);
```

Here:

- \( \text{thesaurus_simple} \) is the new dictionary’s name
- \( \text{mythesaurus} \) is the base name of the thesaurus configuration file. (Its full name will be \( \$\text{SHAREDIR/tsearch_data/mythesaurus.ths} \), where \( \$\text{SHAREDIR} \) means the installation shared-data directory.)
- \( \text{pg_catalog.english_stem} \) is the subdictionary (here, a Snowball English stemmer) to use for thesaurus normalization. Notice that the subdictionary will have its own configuration (for example, stop words), which is not shown here.

Now it is possible to bind the thesaurus dictionary \( \text{thesaurus_simple} \) to the desired token types in a configuration, for example:

```
ALTER TEXT SEARCH CONFIGURATION russian
    ALTER MAPPING FOR asciiword, asciihword, hword_asciipart
        WITH thesaurus_simple;
```
12.6.4.2. Thesaurus Example

Consider a simple astronomical thesaurus `thesaurus_astro`, which contains some astronomical word combinations:

- supernovae stars : sn
- crab nebulae : crab

Below we create a dictionary and bind some token types to an astronomical thesaurus and English stemmer:

```sql
CREATE TEXT SEARCH DICTIONARY thesaurus_astro (    
    TEMPLATE = thesaurus,    
    DictFile = thesaurus_astro,    
    Dictionary = english_stem
);

ALTER TEXT SEARCH CONFIGURATION russian
    
ALTER MAPPING FOR asciiword, asciihword, hword_asciipart
    WITH thesaurus_astro, english_stem;
```

Now we can see how it works. `ts_lexize` is not very useful for testing a thesaurus, because it treats its input as a single token. Instead we can use `plainto_tsquery` and `to_tsvector` which will break their input strings into multiple tokens:

```sql
SELECT plainto_tsquery('supernova star');
plainto_tsquery
-----------------
'sn'

SELECT to_tsvector('supernova star');
to_tsvector
-------------
'sn':1
```

In principle, one can use `to_tsquery` if you quote the argument:

```sql
SELECT to_tsquery('"supernova star"');
to_tsquery
--------------
'sn'
```

Notice that `supernova star` matches `supernovae stars` in `thesaurus_astro` because we specified the `english_stem` stemmer in the thesaurus definition. The stemmer removed the `e` and `s`.

To index the original phrase as well as the substitute, just include it in the right-hand part of the definition:

```sql
supernovae stars : sn supernovae stars

SELECT plainto_tsquery('supernova star');
plainto_tsquery
-------------------------
'sn' & 'supernova' & 'star'
```
12.6.5. Ispell Dictionary

The Ispell dictionary template supports morphological dictionaries, which can normalize many different linguistic forms of a word into the same lexeme. For example, an English Ispell dictionary can match all declensions and conjugations of the search term bank, e.g., banking, banked, banks, banks’, and bank’s.

The standard PostgreSQL distribution does not include any Ispell configuration files. Dictionaries for a large number of languages are available from Ispell\(^1\). Also, some more modern dictionary file formats are supported — MySpell\(^2\) (OO < 2.0.1) and Hunspell\(^3\) (OO >= 2.0.2). A large list of dictionaries is available on the OpenOffice Wiki\(^4\).

To create an Ispell dictionary perform these steps:

- download dictionary configuration files. OpenOffice extension files have the .oxt extension. It is necessary to extract .aff and .dic files, change extensions to .affix and .dict. For some dictionary files it is also needed to convert characters to the UTF-8 encoding with commands (for example, for a Norwegian language dictionary):

  \[
  \text{iconv -f ISO\_8859-1 -t UTF\_8 -o nn\_no.affix nn\_NO.aff} \\
  \text{iconv -f ISO\_8859-1 -t UTF\_8 -o nn\_no.dic nn\_NO.dic}
  \]

- copy files to the $SHAREDIR/tsearch_data directory
- load files into PostgreSQL with the following command:

  \[
  \text{CREATE TEXT SEARCH DICTIONARY english_hunspell (}
  \text{ TEMPLATE = ispell,}
  \text{ DictFile = en\_us,}
  \text{ AffFile = en\_us,}
  \text{ Stopwords = english);}
  \]

Here, DictFile, AffFile, and StopWords specify the base names of the dictionary, affixes, and stop-words files. The stop-words file has the same format explained above for the simple dictionary type. The format of the other files is not specified here but is available from the above-mentioned web sites.

Ispell dictionaries usually recognize a limited set of words, so they should be followed by another broader dictionary; for example, a Snowball dictionary, which recognizes everything.

The .affix file of Ispell has the following structure:

```
prefixes
flag *A:
   .         > RE    # As in enter > reenter
suffixes
flag T:
   E         > ST    # As in late > latest
   [AEIOU]Y  > -Y,IEST # As in dirty > dirtiest
   [AEIOU]Y  > EST    # As in gray > grayest
   [*EY]     > EST    # As in small > smallest
```

And the .dict file has the following structure:

```
lapse/ADGRS
lard/DGRS
```

---

Format of the .dict file is:

`basic_form/affix_class_name`

In the .affix file every affix flag is described in the following format:

`condition > [-stripping_letters,] adding_affix`

Here, condition has a format similar to the format of regular expressions. It can use groupings [..] and[^..]. For example, [AEIOU]Y means that the last letter of the word is "y" and the penultimate letter is "a", "e", "i", "o" or "u". [^EY] means that the last letter is neither "e" nor "y".

Ispell dictionaries support splitting compound words; a useful feature. Notice that the affix file should specify a special flag using the `compoundwords controlled` statement that marks dictionary words that can participate in compound formation:

```
compoundwords controlled z
```

Here are some examples for the Norwegian language:

```
SELECT ts_lexize('norwegian_ispell', 'overbuljongterningpakkmesterassistent');
{over,buljong,terning,pakk,mester,assistent}
SELECT ts_lexize('norwegian_ispell', 'sjokoladefabrikk');
 {sjokoladefabrikk,sjokolade,fabrikk}
```

MySpell format is a subset of Hunspell. The .affix file of Hunspell has the following structure:

```
PFX A Y 1
PFX A 0 re .
SFX T N 4
SFX T 0 st e
SFX T y iest [^aeiou]y
SFX T 0 est [aeiou]y
SFX T 0 est [^ey]
```

The first line of an affix class is the header. Fields of an affix rules are listed after the header:

- parameter name (PFX or SFX)
- flag (name of the affix class)
- stripping characters from beginning (at prefix) or end (at suffix) of the word
- adding affix
- condition that has a format similar to the format of regular expressions.

The .dict file looks like the .dict file of Ispell:

```
larder/M
lardy/RT
```
Chapter 12. Full Text Search

Note: MySpell does not support compound words. Hunspell has sophisticated support for compound words. At present, PostgreSQL implements only the basic compound word operations of Hunspell.

12.6.6. Snowball Dictionary

The Snowball dictionary template is based on a project by Martin Porter, inventor of the popular Porter’s stemming algorithm for the English language. Snowball now provides stemming algorithms for many languages (see the Snowball site\(^5\) for more information). Each algorithm understands how to reduce common variant forms of words to a base, or stem, spelling within its language. A Snowball dictionary requires a language parameter to identify which stemmer to use, and optionally can specify a stopword file name that gives a list of words to eliminate. (PostgreSQL’s standard stopword lists are also provided by the Snowball project.) For example, there is a built-in definition equivalent to

```sql
CREATE TEXT SEARCH DICTIONARY english_stem (
    TEMPLATE = snowball,
    Language = english,
    StopWords = english
);
```

The stopword file format is the same as already explained.

A Snowball dictionary recognizes everything, whether or not it is able to simplify the word, so it should be placed at the end of the dictionary list. It is useless to have it before any other dictionary because a token will never pass through it to the next dictionary.

12.7. Configuration Example

A text search configuration specifies all options necessary to transform a document into a `tsvector`: the parser to use to break text into tokens, and the dictionaries to use to transform each token into a lexeme. Every call of `to_tsvector` or `to_tsquery` needs a text search configuration to perform its processing. The configuration parameter `default_text_search_config` specifies the name of the default configuration, which is the one used by text search functions if an explicit configuration parameter is omitted. It can be set in `postgresql.conf`, or set for an individual session using the `SET` command.

Several predefined text search configurations are available, and you can create custom configurations easily. To facilitate management of text search objects, a set of SQL commands is available, and there are several psql commands that display information about text search objects (Section 12.10).

As an example we will create a configuration `pg`, starting by duplicating the built-in `english` configuration:

```sql
CREATE TEXT SEARCH CONFIGURATION public.pg ( COPY = pg_catalog.english );
```

\(^5\) [http://snowballstem.org/](http://snowballstem.org/)
We will use a PostgreSQL-specific synonym list and store it in $SHAREDIR/tsearch_data/pg_dict.syn. The file contents look like:

postgres pg
pgsql pg
postgresql pg

We define the synonym dictionary like this:

CREATE TEXT SEARCH DICTIONARY pg_dict ( 
   TEMPLATE = synonym, 
   SYNONYMS = pg_dict
);

Next we register the Ispell dictionary english_ispell, which has its own configuration files:

CREATE TEXT SEARCH DICTIONARY english_ispell ( 
   TEMPLATE = ispell, 
   DictFile = english, 
   AffFile = english, 
   StopWords = english
);

Now we can set up the mappings for words in configuration pg:

ALTER TEXT SEARCH CONFIGURATION pg 
   ALTER MAPPING FOR asciiword, asciihword, hword_asciipart, 
      word, hword, hword_part 
      WITH pg_dict, english_ispell, english_stem;

We choose not to index or search some token types that the built-in configuration does handle:

ALTER TEXT SEARCH CONFIGURATION pg 
   DROP MAPPING FOR email, url, url_path, sfloat, float;

Now we can test our configuration:

SELECT * FROM ts_debug('public.pg', 'PostgreSQL, the highly scalable, SQL compliant, open source object-relational database management system, is now undergoing beta testing of the next version of our software.');

The next step is to set the session to use the new configuration, which was created in the public schema:

=> \dF
   List of text search configurations
   Schema | Name | Description
   --------------------------
   public | pg |

SET default_text_search_config = 'public.pg';
SET
12.8. Testing and Debugging Text Search

The behavior of a custom text search configuration can easily become confusing. The functions described in this section are useful for testing text search objects. You can test a complete configuration, or test parsers and dictionaries separately.

12.8.1. Configuration Testing

The function `ts_debug` allows easy testing of a text search configuration.

```sql
ts_debug([ config regconfig, ] document text,
          OUT alias text,
          OUT description text,
          OUT token text,
          OUT dictionaries regdictionary[],
          OUT dictionary regdictionary,
          OUT lexemes text[])
returns setof record
```

`ts_debug` displays information about every token of `document` as produced by the parser and processed by the configured dictionaries. It uses the configuration specified by `config`, or `default_text_search_config` if that argument is omitted.

`ts_debug` returns one row for each token identified in the text by the parser. The columns returned are

- `alias text` — short name of the token type
- `description text` — description of the token type
- `token text` — text of the token
- `dictionaries regdictionary[]` — the dictionaries selected by the configuration for this token type
- `dictionary regdictionary` — the dictionary that recognized the token, or `NULL` if none did
- `lexemes text[]` — the lexeme(s) produced by the dictionary that recognized the token, or `NULL` if none did; an empty array `{}` means it was recognized as a stop word

Here is a simple example:

```sql
SELECT * FROM ts_debug('english','a fat cat sat on a mat - it ate a fat rats');
```

<table>
<thead>
<tr>
<th>alias</th>
<th>description</th>
<th>token</th>
<th>dictionaries</th>
<th>dictionary</th>
<th>lexemes</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>a</td>
<td><code>{english_stem}</code></td>
<td><code>english_stem</code></td>
<td><code>{}</code></td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>fat</td>
<td><code>{english_stem}</code></td>
<td><code>english_stem</code></td>
<td><code>{fat}</code></td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
For a more extensive demonstration, we first create a public.english configuration and Ispell dictionary for the English language:

```
CREATE TEXT SEARCH CONFIGURATION public.english ( COPY = pg_catalog.english );

CREATE TEXT SEARCH DICTIONARY english_ispell {
    TEMPLATE = ispell,
    DictFile = english,
    AffFile = english,
    StopWords = english
};

ALTER TEXT SEARCH CONFIGURATION public.english
ALTER MAPPING FOR asciiword WITH english_ispell, english_stem;
```

```
SELECT * FROM ts_debug('public.english','The Brightest supernovaes');
```

<table>
<thead>
<tr>
<th>alias</th>
<th>description</th>
<th>token</th>
<th>dictionaries</th>
<th>dictionary</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>The</td>
<td>{english_ispell,english_stem}</td>
<td>english_ispell</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td></td>
<td>()</td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>Brightest</td>
<td>{english_ispell,english_stem}</td>
<td>english_ispell</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
<td></td>
<td>()</td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
<td>supernovaes</td>
<td>{english_ispell,english_stem}</td>
<td>english_ispell</td>
</tr>
</tbody>
</table>

In this example, the word `Brightest` was recognized by the parser as an ASCII word (alias `asciiword`). For this token type the dictionary list is `english_ispell` and `english_stem`. The word was recognized by `english_ispell`, which reduced it to the noun `bright`. The word `supernovaes` is unknown to the `english_ispell` dictionary so it was passed to the next dictionary, and, fortunately, was recognized (in fact, `english_stem` is a Snowball dictionary which recognizes everything; that is why it was placed at the end of the dictionary list).

The word `The` was recognized by the `english_ispell` dictionary as a stop word (Section 12.6.1) and will not be indexed. The spaces are discarded too, since the configuration provides no dictionaries at all for them.
You can reduce the width of the output by explicitly specifying which columns you want to see:

```
SELECT alias, token, dictionary, lexemes
FROM ts_debug('public.english','The Brightest supernovae');
```

<table>
<thead>
<tr>
<th>alias</th>
<th>token</th>
<th>dictionary</th>
<th>lexemes</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciiword</td>
<td>The</td>
<td>english_ispell</td>
<td>{}</td>
</tr>
<tr>
<td>blank</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>Brightest</td>
<td>english_ispell</td>
<td>{bright}</td>
</tr>
<tr>
<td>blank</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>asciiword</td>
<td>supernovae</td>
<td>english_stem</td>
<td>{supernova}</td>
</tr>
</tbody>
</table>

12.8.2. Parser Testing

The following functions allow direct testing of a text search parser.

```
ts_parse(parser_name text, document text,
           OUT tokid integer, OUT token text) returns setof record
ts_parse(parser_oid oid, document text,
           OUT tokid integer, OUT token text) returns setof record
```

ts_parse parses the given document and returns a series of records, one for each token produced by parsing. Each record includes a tokid showing the assigned token type and a token which is the text of the token. For example:

```
SELECT * FROM ts_parse('default', '123 - a number');
```

<table>
<thead>
<tr>
<th>tokid</th>
<th>token</th>
</tr>
</thead>
<tbody>
<tr>
<td>22</td>
<td>123</td>
</tr>
<tr>
<td>12</td>
<td>-</td>
</tr>
<tr>
<td>1</td>
<td>a</td>
</tr>
<tr>
<td>12</td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>number</td>
</tr>
</tbody>
</table>

```
ts_token_type(parser_name text, OUT tokid integer,
                 OUT alias text, OUT description text) returns setof record
ts_token_type(parser_oid oid, OUT tokid integer,
                 OUT alias text, OUT description text) returns setof record
```

Ts_token_type returns a table which describes each type of token the specified parser can recognize. For each token type, the table gives the integer tokid that the parser uses to label a token of that type, the alias that names the token type in configuration commands, and a short description. For example:

```
SELECT * FROM ts_token_type('default');
```

<table>
<thead>
<tr>
<th>tokid</th>
<th>alias</th>
<th>description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>asciiword</td>
<td>Word, all ASCII</td>
</tr>
<tr>
<td>2</td>
<td>word</td>
<td>Word, all letters</td>
</tr>
<tr>
<td>3</td>
<td>numword</td>
<td>Word, letters and digits</td>
</tr>
<tr>
<td>4</td>
<td>email</td>
<td>Email address</td>
</tr>
</tbody>
</table>
12.8.3. Dictionary Testing

The `ts_lexize` function facilitates dictionary testing.

\[ \text{ts_lexize(dict regdictionary, token text) returns text[]} \]

`ts_lexize` returns an array of lexemes if the input `token` is known to the dictionary, or an empty array if the token is known to the dictionary but it is a stop word, or `NULL` if it is an unknown word.

Examples:

```
SELECT ts_lexize('english_stem', 'stars');
```

```
|--{star}|
```

```
SELECT ts_lexize('english_stem', 'a');
```

```
|--{}|
```

**Note:** The `ts_lexize` function expects a single `token`, not text. Here is a case where this can be confusing:

```
SELECT ts_lexize('thesaurus_astro', 'supernovae stars') is null;
```

```
?column?
```

```
--t
```

The thesaurus dictionary `thesaurus_astro` does know the phrase `supernovae stars`, but `ts_lexize` fails since it does not parse the input text but treats it as a single token. Use `plainto_tsquery` or `to_tsvector` to test thesaurus dictionaries, for example:
SELECT plainto_tsquery('supernovae stars');
plainto_tsquery
-----------------
'sn'

12.9. GIN and GiST Index Types

There are two kinds of indexes that can be used to speed up full text searches. Note that indexes are not mandatory for full text searching, but in cases where a column is searched on a regular basis, an index is usually desirable.

CREATE INDEX name ON table USING GIN (column);

Creates a GIN (Generalized Inverted Index)-based index. The column must be of tsvector type.

CREATE INDEX name ON table USING GIST (column);

Creates a GiST (Generalized Search Tree)-based index. The column can be of tsvector or tsquery type.

GIN indexes are the preferred text search index type. As inverted indexes, they contain an index entry for each word (lexeme), with a compressed list of matching locations. Multi-word searches can find the first match, then use the index to remove rows that are lacking additional words. GIN indexes store only the words (lexemes) of tsvector values, and not their weight labels. Thus a table row recheck is needed when using a query that involves weights.

A GiST index is lossy, meaning that the index might produce false matches, and it is necessary to check the actual table row to eliminate such false matches. (PostgreSQL does this automatically when needed.) GiST indexes are lossy because each document is represented in the index by a fixed-length signature. The signature is generated by hashing each word into a single bit in an n-bit string, with all these bits OR-ed together to produce an n-bit document signature. When two words hash to the same bit position there will be a false match. If all words in the query have matches (real or false) then the table row must be retrieved to see if the match is correct.

Lossiness causes performance degradation due to unnecessary fetches of table records that turn out to be false matches. Since random access to table records is slow, this limits the usefulness of GiST indexes. The likelihood of false matches depends on several factors, in particular the number of unique words, so using dictionaries to reduce this number is recommended.

Note that GIN index build time can often be improved by increasing maintenance_work_mem, while GiST index build time is not sensitive to that parameter.

Partitioning of big collections and the proper use of GIN and GiST indexes allows the implementation of very fast searches with online update. Partitioning can be done at the database level using table inheritance, or by distributing documents over servers and collecting search results using the dblink module. The latter is possible because ranking functions use only local information.
12.10. psql Support

Information about text search configuration objects can be obtained in psql using a set of commands:

```
\dF{d,p,t}[+] [PATTERN]
```

An optional + produces more details.

The optional parameter PATTERN can be the name of a text search object, optionally schema-qualified. If PATTERN is omitted then information about all visible objects will be displayed. PATTERN can be a regular expression and can provide separate patterns for the schema and object names. The following examples illustrate this:

```
=> \dF *fulltext*
   List of text search configurations
   Schema | Name   | Description
       |--------+-------------------
   public | fulltext_cfg |

=> \dF *.fulltext*
   List of text search configurations
   Schema | Name   | Description
       |--------+-------------------
   fulltext | fulltext_cfg |
   public  | fulltext_cfg |
```

The available commands are:

```
\dF{+} [PATTERN]
   List text search configurations (add + for more detail).
```

```
=> \dF russian
   List of text search configurations
   Schema | Name   | Description
       |--------+-------------------
   pg_catalog | russian | configuration for russian language

=> \dF+ russian
Text search configuration "pg_catalog.russian"
Parser: "pg_catalog.default"
   Token         | Dictionaries
----------------+--------------
   asciihword     | english.stem
   asciiword      | english.stem
   email          | simple
   file           | simple
   float          | simple
   host           | simple
   hword          | russian.stem
   hword_asciipart | english.stem
   hword_numpart  | simple
   hword_part     | russian.stem
   int            | simple
   numhword       | simple
   numword        | simple
   sfloat         | simple
   uint           | simple
   url            | simple
```

411
List text search dictionaries (add + for more detail).

=> \dFd+

<table>
<thead>
<tr>
<th>Schema</th>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_catalog</td>
<td>danish_stem</td>
<td>snowball stemmer for danish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>dutch_stem</td>
<td>snowball stemmer for dutch language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>english_stem</td>
<td>snowball stemmer for english language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>finnish_stem</td>
<td>snowball stemmer for finnish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>french_stem</td>
<td>snowball stemmer for french language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>german_stem</td>
<td>snowball stemmer for german language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>hungarian_stem</td>
<td>snowball stemmer for hungarian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>italian_stem</td>
<td>snowball stemmer for italian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>norwegian_stem</td>
<td>snowball stemmer for norwegian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>portuguese_stem</td>
<td>snowball stemmer for portuguese language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>romanian_stem</td>
<td>snowball stemmer for romanian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>russian_stem</td>
<td>snowball stemmer for russian language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>simple</td>
<td>simple dictionary: just lower case and check for stop words</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>spanish_stem</td>
<td>snowball stemmer for spanish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>swedish_stem</td>
<td>snowball stemmer for swedish language</td>
</tr>
<tr>
<td>pg_catalog</td>
<td>turkish_stem</td>
<td>snowball stemmer for turkish language</td>
</tr>
</tbody>
</table>

List text search parsers (add + for more detail).

=> \dFp+

<table>
<thead>
<tr>
<th>Schema</th>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_catalog</td>
<td>default</td>
<td>default word parser</td>
</tr>
</tbody>
</table>

=> \dFp+

Text search parser "pg_catalog.default"

<table>
<thead>
<tr>
<th>Method</th>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Start parse</td>
<td>prsd_start</td>
<td></td>
</tr>
<tr>
<td>Get next token</td>
<td>prsd_nexttoken</td>
<td></td>
</tr>
<tr>
<td>End parse</td>
<td>prsd_end</td>
<td></td>
</tr>
<tr>
<td>Get headline</td>
<td>prsd_headline</td>
<td></td>
</tr>
<tr>
<td>Get token types</td>
<td>prsd_lextype</td>
<td></td>
</tr>
</tbody>
</table>

Token types for parser "pg_catalog.default"

<table>
<thead>
<tr>
<th>Token name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>asciihword</td>
<td>Hyphenated word, all ASCII</td>
</tr>
<tr>
<td>asciiword</td>
<td>Word, all ASCII</td>
</tr>
<tr>
<td>blank</td>
<td>Space symbols</td>
</tr>
<tr>
<td>email</td>
<td>Email address</td>
</tr>
<tr>
<td>entity</td>
<td>XML entity</td>
</tr>
<tr>
<td>file</td>
<td>File or path name</td>
</tr>
<tr>
<td>float</td>
<td>Decimal notation</td>
</tr>
<tr>
<td>host</td>
<td>Host</td>
</tr>
<tr>
<td>hword</td>
<td>Hyphenated word, all letters</td>
</tr>
</tbody>
</table>
12.11. Limitations

The current limitations of PostgreSQL’s text search features are:

- The length of each lexeme must be less than 2K bytes
- The length of a `tsvector` (lexemes + positions) must be less than 1 megabyte
- The number of lexemes must be less than 2^64
- Position values in `tsvector` must be greater than 0 and no more than 16,383
- The match distance in a `<N>` (FOLLOWED BY) `tsquery` operator cannot be more than 16,384
- No more than 256 positions per lexeme
- The number of nodes (lexemes + operators) in a `tsquery` must be less than 32,768

For comparison, the PostgreSQL 8.1 documentation contained 10,441 unique words, a total of 335,420 words, and the most frequent word “postgresql” was mentioned 6,127 times in 655 documents.

Another example — the PostgreSQL mailing list archives contained 910,989 unique words with 57,491,343 lexemes in 461,020 messages.
12.12. Migration from Pre-8.3 Text Search

Applications that use the tsearch2 module for text searching will need some adjustments to work with the built-in features:

- Some functions have been renamed or had small adjustments in their argument lists, and all of them are now in the `pg_catalog` schema, whereas in a previous installation they would have been in `public` or another non-system schema. There is a new version of tsearch2 that provides a compatibility layer to solve most problems in this area.

- The old tsearch2 functions and other objects must be suppressed when loading pg_dump output from a pre-8.3 database. While many of them won’t load anyway, a few will and then cause problems. One simple way to deal with this is to load the new tsearch2 module before restoring the dump; then it will block the old objects from being loaded.

- Text search configuration setup is completely different now. Instead of manually inserting rows into configuration tables, search is configured through the specialized SQL commands shown earlier in this chapter. There is no automated support for converting an existing custom configuration for 8.3; you’re on your own here.

- Most types of dictionaries rely on some outside-the-database configuration files. These are largely compatible with pre-8.3 usage, but note the following differences:
  - Configuration files now must be placed in a single specified directory (`$SHAREDIR/tsearch_data`), and must have a specific extension depending on the type of file, as noted previously in the descriptions of the various dictionary types. This restriction was added to forestall security problems.
  - Configuration files must be encoded in UTF-8 encoding, regardless of what database encoding is used.
  - In thesaurus configuration files, stop words must be marked with `?`.
Chapter 13. Concurrency Control

This chapter describes the behavior of the PostgreSQL database system when two or more sessions try to access the same data at the same time. The goals in that situation are to allow efficient access for all sessions while maintaining strict data integrity. Every developer of database applications should be familiar with the topics covered in this chapter.

13.1. Introduction

PostgreSQL provides a rich set of tools for developers to manage concurrent access to data. Internally, data consistency is maintained by using a multiversion model (Multiversion Concurrency Control, MVCC). This means that each SQL statement sees a snapshot of data (a database version) as it was some time ago, regardless of the current state of the underlying data. This prevents statements from viewing inconsistent data produced by concurrent transactions performing updates on the same data rows, providing transaction isolation for each database session. MVCC, by eschewing the locking methodologies of traditional database systems, minimizes lock contention in order to allow for reasonable performance in multiuser environments.

The main advantage of using the MVCC model of concurrency control rather than locking is that in MVCC locks acquired for querying (reading) data do not conflict with locks acquired for writing data, and so reading never blocks writing and writing never blocks reading. PostgreSQL maintains this guarantee even when providing the strictest level of transaction isolation through the use of an innovative Serializable Snapshot Isolation (SSI) level.

Table- and row-level locking facilities are also available in PostgreSQL for applications which don’t generally need full transaction isolation and prefer to explicitly manage particular points of conflict. However, proper use of MVCC will generally provide better performance than locks. In addition, application-defined advisory locks provide a mechanism for acquiring locks that are not tied to a single transaction.

13.2. Transaction Isolation

The SQL standard defines four levels of transaction isolation. The most strict is Serializable, which is defined by the standard in a paragraph which says that any concurrent execution of a set of Serializable transactions is guaranteed to produce the same effect as running them one at a time in some order. The other three levels are defined in terms of phenomena, resulting from interaction between concurrent transactions, which must not occur at each level. The standard notes that due to the definition of Serializable, none of these phenomena are possible at that level. (This is hardly surprising -- if the effect of the transactions must be consistent with having been run one at a time, how could you see any phenomena caused by interactions?)

The phenomena which are prohibited at various levels are:

dirty read

A transaction reads data written by a concurrent uncommitted transaction.

nonrepeatable read

A transaction re-reads data it has previously read and finds that data has been modified by another transaction (that committed since the initial read).
phantom read

A transaction re-executes a query returning a set of rows that satisfy a search condition and finds that the set of rows satisfying the condition has changed due to another recently-committed transaction.

serialization anomaly

The result of successfully committing a group of transactions is inconsistent with all possible orderings of running those transactions one at a time.

The SQL standard and PostgreSQL-implemented transaction isolation levels are described in Table 13-1.

Table 13-1. Transaction Isolation Levels

<table>
<thead>
<tr>
<th>Isolation Level</th>
<th>Dirty Read</th>
<th>Nonrepeatable Read</th>
<th>Phantom Read</th>
<th>Serialization Anomaly</th>
</tr>
</thead>
<tbody>
<tr>
<td>Read uncommitted</td>
<td>Allowed, but not</td>
<td>Possible</td>
<td>Possible</td>
<td>Possible</td>
</tr>
<tr>
<td></td>
<td>in PG</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Read committed</td>
<td>Not possible</td>
<td>Possible</td>
<td>Possible</td>
<td>Possible</td>
</tr>
<tr>
<td>Repeatable read</td>
<td>Not possible</td>
<td>Not possible</td>
<td>Allowed, but not</td>
<td>Possible</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>in PG</td>
<td></td>
</tr>
<tr>
<td>Serializable</td>
<td>Not possible</td>
<td>Not possible</td>
<td>Not possible</td>
<td>Not possible</td>
</tr>
</tbody>
</table>

In PostgreSQL, you can request any of the four standard transaction isolation levels, but internally only three distinct isolation levels are implemented, i.e. PostgreSQL’s Read Uncommitted mode behaves like Read Committed. This is because it is the only sensible way to map the standard isolation levels to PostgreSQL’s multiversion concurrency control architecture.

The table also shows that PostgreSQL’s Repeatable Read implementation does not allow phantom reads. Stricter behavior is permitted by the SQL standard: the four isolation levels only define which phenomena must not happen, not which phenomena must happen. The behavior of the available isolation levels is detailed in the following subsections.

To set the transaction isolation level of a transaction, use the command SET TRANSACTION.

**Important:** Some PostgreSQL data types and functions have special rules regarding transactional behavior. In particular, changes made to a sequence (and therefore the counter of a column declared using `serial`) are immediately visible to all other transactions and are not rolled back if the transaction that made the changes aborts. See Section 9.16 and Section 8.1.4.

### 13.2.1. Read Committed Isolation Level

*Read Committed* is the default isolation level in PostgreSQL. When a transaction uses this isolation level, a `SELECT` query (without a `FOR UPDATE/SHARE` clause) sees only data committed before the query began; it never sees either uncommitted data or changes committed during query execution by concurrent transactions. In effect, a `SELECT` query sees a snapshot of the database as of the instant the query begins to run. However, `SELECT` does see the effects of previous updates executed within its own transaction, even though they are not yet committed. Also note that two successive `SELECT` commands can see different data, even though they are within a single transaction, if other transactions commit changes after the first `SELECT` starts and before the second `SELECT` starts.
Chapter 13. Concurrency Control

UPDATE, DELETE, SELECT FOR UPDATE, and SELECT FOR SHARE commands behave the same as SELECT in terms of searching for target rows: they will only find target rows that were committed as of the command start time. However, such a target row might have already been updated (or deleted or locked) by another concurrent transaction by the time it is found. In this case, the would-be updater will wait for the first updating transaction to commit or roll back (if it is still in progress). If the first updater rolls back, then its effects are negated and the second updater can proceed with updating the originally found row. If the first updater commits, the second updater will ignore the row if the first updater deleted it, otherwise it will attempt to apply its operation to the updated version of the row. The search condition of the command (the WHERE clause) is re-evaluated to see if the updated version of the row still matches the search condition. If so, the second updater proceeds with its operation using the updated version of the row. In the case of SELECT FOR UPDATE and SELECT FOR SHARE, this means it is the updated version of the row that is locked and returned to the client.

INSERT with an ON CONFLICT DO UPDATE clause behaves similarly. In Read Committed mode, each row proposed for insertion will either insert or update. Unless there are unrelated errors, one of those two outcomes is guaranteed. If a conflict originates in another transaction whose effects are not yet visible to the INSERT, the UPDATE clause will affect that row, even though possibly no version of that row is conventionally visible to the command.

INSERT with an ON CONFLICT DO NOTHING clause may have insertion not proceed for a row due to the outcome of another transaction whose effects are not visible to the INSERT snapshot. Again, this is only the case in Read Committed mode.

Because of the above rules, it is possible for an updating command to see an inconsistent snapshot: it can see the effects of concurrent updating commands on the same rows it is trying to update, but it does not see effects of those commands on other rows in the database. This behavior makes Read Committed mode unsuitable for commands that involve complex search conditions; however, it is just right for simpler cases. For example, consider updating bank balances with transactions like:

BEGIN;
UPDATE accounts SET balance = balance + 100.00 WHERE acctnum = 12345;
UPDATE accounts SET balance = balance - 100.00 WHERE acctnum = 7534;
COMMIT;

If two such transactions concurrently try to change the balance of account 12345, we clearly want the second transaction to start with the updated version of the account’s row. Because each command is affecting only a predetermined row, letting it see the updated version of the row does not create any troublesome inconsistency.

More complex usage can produce undesirable results in Read Committed mode. For example, consider a DELETE command operating on data that is being both added and removed from its restriction criteria by another command, e.g., assume website is a two-row table with website.hits equaling 9 and 10:

BEGIN;
UPDATE website SET hits = hits + 1;
-- run from another session: DELETE FROM website WHERE hits = 10;
COMMIT;

The DELETE will have no effect even though there is a website.hits = 10 row before and after the UPDATE. This occurs because the pre-update row value 9 is skipped, and when the UPDATE completes and DELETE obtains a lock, the new row value is no longer 10 but 11, which no longer matches the criteria.

Because Read Committed mode starts each command with a new snapshot that includes all transactions committed up to that instant, subsequent commands in the same transaction will see the effects
Chapter 13. Concurrency Control

of the committed concurrent transaction in any case. The point at issue above is whether or not a single command sees an absolutely consistent view of the database.

The partial transaction isolation provided by Read Committed mode is adequate for many applications, and this mode is fast and simple to use; however, it is not sufficient for all cases. Applications that do complex queries and updates might require a more rigorously consistent view of the database than Read Committed mode provides.

13.2.2. Repeatable Read Isolation Level

The Repeatable Read isolation level only sees data committed before the transaction began; it never sees either uncommitted data or changes committed during transaction execution by concurrent transactions. (However, the query does see the effects of previous updates executed within its own transaction, even though they are not yet committed.) This is a stronger guarantee than is required by the SQL standard for this isolation level, and prevents all of the phenomena described in Table 13-1 except for serialization anomalies. As mentioned above, this is specifically allowed by the standard, which only describes the minimum protections each isolation level must provide.

This level is different from Read Committed in that a query in a repeatable read transaction sees a snapshot as of the start of the first non-transaction-control statement in the transaction, not as of the start of the current statement within the transaction. Thus, successive SELECT commands within a single transaction see the same data, i.e., they do not see changes made by other transactions that committed after their own transaction started.

Applications using this level must be prepared to retry transactions due to serialization failures. UPDATE, DELETE, SELECT FOR UPDATE, and SELECT FOR SHARE commands behave the same as SELECT in terms of searching for target rows: they will only find target rows that were committed as of the transaction start time. However, such a target row might have already been updated (or deleted or locked) by another concurrent transaction by the time it is found. In this case, the repeatable read transaction will wait for the first updating transaction to commit or roll back (if it is still in progress). If the first updater rolls back, then its effects are negated and the repeatable read transaction can proceed with updating the originally found row. But if the first updater commits (and actually updated or deleted the row, not just locked it) then the repeatable read transaction will be rolled back with the message

ERROR: could not serialize access due to concurrent update

because a repeatable read transaction cannot modify or lock rows changed by other transactions after the repeatable read transaction began.

When an application receives this error message, it should abort the current transaction and retry the whole transaction from the beginning. The second time through, the transaction will see the previously-committed change as part of its initial view of the database, so there is no logical conflict in using the new version of the row as the starting point for the new transaction’s update.

Note that only updating transactions might need to be retried; read-only transactions will never have serialization conflicts.

The Repeatable Read mode provides a rigorous guarantee that each transaction sees a completely stable view of the database. However, this view will not necessarily always be consistent with some serial (one at a time) execution of concurrent transactions of the same level. For example, even a read only transaction at this level may see a control record updated to show that a batch has been completed but not see one of the detail records which is logically part of the batch because it read an earlier revision of the control record. Attempts to enforce business rules by transactions running
at this isolation level are not likely to work correctly without careful use of explicit locks to block conflicting transactions.

**Note:** Prior to PostgreSQL version 9.1, a request for the Serializable transaction isolation level provided exactly the same behavior described here. To retain the legacy Serializable behavior, Repeatable Read should now be requested.

### 13.2.3. Serializable Isolation Level

The *Serializable* isolation level provides the strictest transaction isolation. This level emulates serial transaction execution for all committed transactions; as if transactions had been executed one after another, serially, rather than concurrently. However, like the Repeatable Read level, applications using this level must be prepared to retry transactions due to serialization failures. In fact, this isolation level works exactly the same as Repeatable Read except that it monitors for conditions which could make execution of a concurrent set of serializable transactions behave in a manner inconsistent with all possible serial (one at a time) executions of those transactions. This monitoring does not introduce any blocking beyond that present in repeatable read, but there is some overhead to the monitoring, and detection of the conditions which could cause a serialization anomaly will trigger a serialization failure.

As an example, consider a table `mytab`, initially containing:

<table>
<thead>
<tr>
<th>class</th>
<th>value</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>10</td>
</tr>
<tr>
<td>1</td>
<td>20</td>
</tr>
<tr>
<td>2</td>
<td>100</td>
</tr>
<tr>
<td>2</td>
<td>200</td>
</tr>
</tbody>
</table>

Suppose that serializable transaction A computes:

```sql
SELECT SUM(value) FROM mytab WHERE class = 1;
```

and then inserts the result (30) as the value in a new row with `class = 2`. Concurrently, serializable transaction B computes:

```sql
SELECT SUM(value) FROM mytab WHERE class = 2;
```

and obtains the result 300, which it inserts in a new row with `class = 1`. Then both transactions try to commit. If either transaction were running at the Repeatable Read isolation level, both would be allowed to commit; but since there is no serial order of execution consistent with the result, using Serializable transactions will allow one transaction to commit and will roll the other back with this message:

**ERROR:** could not serialize access due to read/write dependencies among transactions

This is because if A had executed before B, B would have computed the sum 330, not 300, and similarly the other order would have resulted in a different sum computed by A.

When relying on Serializable transactions to prevent anomalies, it is important that any data read from a permanent user table not be considered valid until the transaction which read it has successfully committed. This is true even for read-only transactions, except that data read within a *deferrable* read-only transaction is known to be valid as soon as it is read, because such a transaction waits until
it can acquire a snapshot guaranteed to be free from such problems before starting to read any data. In all other cases applications must not depend on results read during a transaction that later aborted; instead, they should retry the transaction until it succeeds.

To guarantee true serializability PostgreSQL uses predicate locking, which means that it keeps locks which allow it to determine when a write would have had an impact on the result of a previous read from a concurrent transaction, had it run first. In PostgreSQL these locks do not cause any blocking and therefore can not play any part in causing a deadlock. They are used to identify and flag dependencies among concurrent Serializable transactions which in certain combinations can lead to serialization anomalies. In contrast, a Read Committed or Repeatable Read transaction which wants to ensure data consistency may need to take out a lock on an entire table, which could block other users attempting to use that table, or it may use `SELECT FOR UPDATE` or `SELECT FOR SHARE` which not only can block other transactions but cause disk access.

Predicate locks in PostgreSQL, like in most other database systems, are based on data actually accessed by a transaction. These will show up in the `pg_locks` system view with a mode of `SIReadLock`. The particular locks acquired during execution of a query will depend on the plan used by the query, and multiple finer-grained locks (e.g., tuple locks) may be combined into fewer coarser-grained locks (e.g., page locks) during the course of the transaction to prevent exhaustion of the memory used to track the locks. A `READ ONLY` transaction may be able to release its SIRead locks before completion, if it detects that no conflicts can still occur which could lead to a serialization anomaly. In fact, `READ ONLY` transactions will often be able to establish that fact at startup and avoid taking any predicate locks. If you explicitly request a `SERIALIZABLE READ ONLY DEFERRABLE` transaction, it will block until it can establish this fact. (This is the only case where Serializable transactions block but Repeatable Read transactions don’t.) On the other hand, SIRead locks often need to be kept past transaction commit, until overlapping read write transactions complete.

Consistent use of Serializable transactions can simplify development. The guarantee that any set of successfully committed concurrent Serializable transactions will have the same effect as if they were run one at a time means that if you can demonstrate that a single transaction, as written, will do the right thing when run by itself, you can have confidence that it will do the right thing in any mix of Serializable transactions, even without any information about what those other transactions might do, or it will not successfully commit. It is important that an environment which uses this technique have a generalized way of handling serialization failures (which always return with a SQLSTATE value of ‘40001’), because it will be very hard to predict exactly which transactions might contribute to the read/write dependencies and need to be rolled back to prevent serialization anomalies. The monitoring of read/write dependencies has a cost, as does the restart of transactions which are terminated with a serialization failure, but balanced against the cost and blocking involved in use of explicit locks and `SELECT FOR UPDATE` or `SELECT FOR SHARE`, Serializable transactions are the best performance choice for some environments.

While PostgreSQL’s Serializable transaction isolation level only allows concurrent transactions to commit if it can prove there is a serial order of execution that would produce the same effect, it doesn’t always prevent errors from being raised that would not occur in true serial execution. In particular, it is possible to see unique constraint violations caused by conflicts with overlapping Serializable transactions even after explicitly checking that the key isn’t present before attempting to insert it. This can be avoided by making sure that all Serializable transactions that insert potentially conflicting keys explicitly check if they can do so first. For example, imagine an application that asks the user for a new key and then checks that it doesn’t exist already by trying to select it first, or generates a new key by selecting the maximum existing key and adding one. If some Serializable transactions insert new keys directly without following this protocol, unique constraints violations might be reported even in cases where they could not occur in a serial execution of the concurrent transactions.

For optimal performance when relying on Serializable transactions for concurrency control, these
issues should be considered:

- Declare transactions as `READ ONLY` when possible.
- Control the number of active connections, using a connection pool if needed. This is always an important performance consideration, but it can be particularly important in a busy system using Serializable transactions.
- Don’t put more into a single transaction than needed for integrity purposes.
- Don’t leave connections dangling “idle in transaction” longer than necessary. The configuration parameter `idle_in_transaction_session_timeout` may be used to automatically disconnect lingering sessions.
- Eliminate explicit locks, `SELECT FOR UPDATE`, and `SELECT FOR SHARE` where no longer needed due to the protections automatically provided by Serializable transactions.
- When the system is forced to combine multiple page-level predicate locks into a single relation-level predicate lock because the predicate lock table is short of memory, an increase in the rate of serialization failures may occur. You can avoid this by increasing `max_pred_locks_per_transaction`.
- A sequential scan will always necessitate a relation-level predicate lock. This can result in an increased rate of serialization failures. It may be helpful to encourage the use of index scans by reducing `random_page_cost` and/or increasing `cpu_tuple_cost`. Be sure to weigh any decrease in transaction rollbacks and restarts against any overall change in query execution time.

### 13.3. Explicit Locking

PostgreSQL provides various lock modes to control concurrent access to data in tables. These modes can be used for application-controlled locking in situations where MVCC does not give the desired behavior. Also, most PostgreSQL commands automatically acquire locks of appropriate modes to ensure that referenced tables are not dropped or modified in incompatible ways while the command executes. (For example, `TRUNCATE` cannot safely be executed concurrently with other operations on the same table, so it obtains an exclusive lock on the table to enforce that.)

To examine a list of the currently outstanding locks in a database server, use the `pg_locks` system view. For more information on monitoring the status of the lock manager subsystem, refer to Chapter 28.

#### 13.3.1. Table-level Locks

The list below shows the available lock modes and the contexts in which they are used automatically by PostgreSQL. You can also acquire any of these locks explicitly with the command `LOCK`. Remember that all of these lock modes are table-level locks, even if the name contains the word “row”; the names of the lock modes are historical. To some extent the names reflect the typical usage of each lock mode — but the semantics are all the same. The only real difference between one lock mode and another is the set of lock modes with which each conflicts (see Table 13-2). Two transactions cannot hold locks of conflicting modes on the same table at the same time. (However, a transaction never conflicts with itself. For example, it might acquire `ACCESS EXCLUSIVE` lock and later acquire `ACCESS SHARE` lock on the same table.) Non-conflicting lock modes can be held concurrently by many transactions. Notice in particular that some lock modes are self-conflicting (for example, an
ACCESS EXCLUSIVE lock cannot be held by more than one transaction at a time) while others are not self-conflicting (for example, an ACCESS SHARE lock can be held by multiple transactions).

**Table-level Lock Modes**

**ACCESS SHARE**
- Conflicts with the ACCESS EXCLUSIVE lock mode only.
- The SELECT command acquires a lock of this mode on referenced tables. In general, any query that only reads a table and does not modify it will acquire this lock mode.

**ROW SHARE**
- Conflicts with the EXCLUSIVE and ACCESS EXCLUSIVE lock modes.
- The SELECT FOR UPDATE and SELECT FOR SHARE commands acquire a lock of this mode on the target table(s) (in addition to ACCESS SHARE locks on any other tables that are referenced but not selected FOR UPDATE/FOR SHARE).

**ROW EXCLUSIVE**
- Conflicts with the SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes.
- The commands UPDATE, DELETE, and INSERT acquire this lock mode on the target table (in addition to ACCESS SHARE locks on any other referenced tables). In general, this lock mode will be acquired by any command that modifies data in a table.

**SHARE UPDATE EXCLUSIVE**
- Conflicts with the SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode protects a table against concurrent schema changes and VACUUM runs.
- Acquired by VACUUM (without FULL), ANALYZE, CREATE INDEX CONCURRENTLY, and ALTER TABLE VALIDATE and other ALTER TABLE variants (for full details see ALTER TABLE).

**SHARE**
- Conflicts with the ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode protects a table against concurrent data changes.
- Acquired by CREATE INDEX (without CONCURRENTLY).

**SHARE ROW EXCLUSIVE**
- Conflicts with the ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode protects a table against concurrent data changes, and is self-exclusive so that only one session can hold it at a time.
- Acquired by CREATE TRIGGER and many forms of ALTER TABLE (see ALTER TABLE).

**EXCLUSIVE**
- Conflicts with the ROW SHARE, ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE lock modes. This mode allows only concurrent ACCESS SHARE locks, i.e., only reads from the table can proceed in parallel with a transaction holding this lock mode.
- Acquired by REFRESH MATERIALIZED VIEW CONCURRENTLY.
Chapter 13. Concurrency Control

ACCESS EXCLUSIVE

Conflicts with locks of all modes (ACCESS SHARE, ROW SHARE, ROW EXCLUSIVE, SHARE UPDATE EXCLUSIVE, SHARE, SHARE ROW EXCLUSIVE, EXCLUSIVE, and ACCESS EXCLUSIVE). This mode guarantees that the holder is the only transaction accessing the table in any way.

Acquired by the DROP TABLE, TRUNCATE, REINDEX, CLUSTER, VACUUM FULL, and REFRESH MATERIALIZED VIEW (without CONCURRENTLY) commands. Many forms of ALTER TABLE also acquire a lock at this level. This is also the default lock mode for LOCK TABLE statements that do not specify a mode explicitly.

Tip: Only an ACCESS EXCLUSIVE lock blocks a SELECT (without FOR UPDATE/SHARE) statement.

Once acquired, a lock is normally held till end of transaction. But if a lock is acquired after establishing a savepoint, the lock is released immediately if the savepoint is rolled back to. This is consistent with the principle that ROLLBACK cancels all effects of the commands since the savepoint. The same holds for locks acquired within a PL/pgSQL exception block: an error escape from the block releases locks acquired within it.

Table 13-2. Conflicting Lock Modes

<table>
<thead>
<tr>
<th>Requested Lock Mode</th>
<th>ACCESS SHARE</th>
<th>ROW SHARE</th>
<th>ROW EXCLUSIVE</th>
<th>SHARE UPDATE EXCLUSIVE</th>
<th>SHARE ROW EXCLUSIVE</th>
<th>EXCLUSIVE</th>
<th>ACCESS EXCLUSIVE</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACCESS SHARE</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>X</td>
</tr>
<tr>
<td>ROW SHARE</td>
<td></td>
<td>X</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ROW EXCLUSIVE</td>
<td></td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td></td>
</tr>
<tr>
<td>SHARE UPDATE EXCLUSIVE</td>
<td></td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td></td>
</tr>
<tr>
<td>SHARE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td></td>
</tr>
<tr>
<td>SHARE ROW EXCLUSIVE</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### 13.3.2. Row-level Locks

In addition to table-level locks, there are row-level locks, which are listed as below with the contexts in which they are used automatically by PostgreSQL. See Table 13-3 for a complete table of row-level lock conflicts. Note that a transaction can hold conflicting locks on the same row, even in different subtransactions; but other than that, two transactions can never hold conflicting locks on the same row. Row-level locks do not affect data querying; they block only *writers and lockers* to the same row.

#### Row-level Lock Modes

<table>
<thead>
<tr>
<th>Requested Lock Mode</th>
<th>Current Lock Mode</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>ACCESS SHARE</td>
</tr>
<tr>
<td>EXCLUSIVE</td>
<td>X</td>
</tr>
<tr>
<td>ACCESS EXCLUSIVE</td>
<td>X</td>
</tr>
</tbody>
</table>

**FOR UPDATE**

*FOR UPDATE* causes the rows retrieved by the *SELECT* statement to be locked as though for update. This prevents them from being locked, modified or deleted by other transactions until the current transaction ends. That is, other transactions that attempt *UPDATE*, *DELETE*, *SELECT FOR UPDATE*, *SELECT FOR NO KEY UPDATE*, *SELECT FOR SHARE* or *SELECT FOR KEY SHARE* of these rows will be blocked until the current transaction ends; conversely, *SELECT FOR UPDATE* will wait for a concurrent transaction that has run any of those commands on the same row, and will then lock and return the updated row (or no row, if the row was deleted). Within a *REPEATABLE READ* or *SERIALIZABLE* transaction, however, an error will be thrown if a row to be locked has changed since the transaction started. For further discussion see Section 13.4.

The *FOR UPDATE* lock mode is also acquired by any *DELETE* on a row, and also by an *UPDATE* that modifies the values on certain columns. Currently, the set of columns considered for the *UPDATE* case are those that have a unique index on them that can be used in a foreign key (so partial indexes and expressional indexes are not considered), but this may change in the future.

**FOR NO KEY UPDATE**

Behaves similarly to *FOR UPDATE*, except that the lock acquired is weaker: this lock will not block *SELECT FOR KEY SHARE* commands that attempt to acquire a lock on the same rows. This lock mode is also acquired by any *UPDATE* that does not acquire a *FOR UPDATE* lock.

**FOR SHARE**

Behaves similarly to *FOR NO KEY UPDATE*, except that it acquires a shared lock rather than exclusive lock on each retrieved row. A shared lock blocks other transactions from performing...
UPDATE, DELETE, SELECT FOR UPDATE or SELECT FOR NO KEY UPDATE on these rows, but it does not prevent them from performing SELECT FOR SHARE or SELECT FOR KEY SHARE.

**FOR KEY SHARE**

Behaves similarly to FOR SHARE, except that the lock is weaker: SELECT FOR UPDATE is blocked, but not SELECT FOR NO KEY UPDATE. A key-shared lock blocks other transactions from performing DELETE or any UPDATE that changes the key values, but not other UPDATE, and neither does it prevent SELECT FOR NO KEY UPDATE, SELECT FOR SHARE, or SELECT FOR KEY SHARE.

PostgreSQL doesn’t remember any information about modified rows in memory, so there is no limit on the number of rows locked at one time. However, locking a row might cause a disk write, e.g., SELECT FOR UPDATE modifies selected rows to mark them locked, and so will result in disk writes.

**Table 13-3. Conflicting Row-level Locks**

<table>
<thead>
<tr>
<th>Requested Lock Mode</th>
<th>Current Lock Mode</th>
<th>FOR KEY SHARE</th>
<th>FOR SHARE</th>
<th>FOR NO KEY UPDATE</th>
<th>FOR UPDATE</th>
</tr>
</thead>
<tbody>
<tr>
<td>FOR KEY SHARE</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>X</td>
</tr>
<tr>
<td>FOR SHARE</td>
<td></td>
<td></td>
<td></td>
<td>X</td>
<td>X</td>
</tr>
<tr>
<td>FOR NO KEY UPDATE</td>
<td></td>
<td>X</td>
<td></td>
<td>X</td>
<td>X</td>
</tr>
<tr>
<td>FOR UPDATE</td>
<td></td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>X</td>
</tr>
</tbody>
</table>

**13.3.3. Page-level Locks**

In addition to table and row locks, page-level share/exclusive locks are used to control read/write access to table pages in the shared buffer pool. These locks are released immediately after a row is fetched or updated. Application developers normally need not be concerned with page-level locks, but they are mentioned here for completeness.

**13.3.4. Deadlocks**

The use of explicit locking can increase the likelihood of deadlocks, wherein two (or more) transactions each hold locks that the other wants. For example, if transaction 1 acquires an exclusive lock on table A and then tries to acquire an exclusive lock on table B, while transaction 2 has already exclusive-locked table B and now wants an exclusive lock on table A, then neither one can proceed. PostgreSQL automatically detects deadlock situations and resolves them by aborting one of the transactions involved, allowing the other(s) to complete. (Exactly which transaction will be aborted is difficult to predict and should not be relied upon.)

Note that deadlocks can also occur as the result of row-level locks (and thus, they can occur even if explicit locking is not used). Consider the case in which two concurrent transactions modify a table. The first transaction executes:

```sql
UPDATE accounts SET balance = balance + 100.00 WHERE acctnum = 11111;
```
This acquires a row-level lock on the row with the specified account number. Then, the second transaction executes:

```sql
UPDATE accounts SET balance = balance + 100.00 WHERE acctnum = 22222;
UPDATE accounts SET balance = balance - 100.00 WHERE acctnum = 11111;
```

The first `UPDATE` statement successfully acquires a row-level lock on the specified row, so it succeeds in updating that row. However, the second `UPDATE` statement finds that the row it is attempting to update has already been locked, so it waits for the transaction that acquired the lock to complete. Transaction two is now waiting on transaction one to complete before it continues execution. Now, transaction one executes:

```sql
UPDATE accounts SET balance = balance - 100.00 WHERE acctnum = 22222;
```

Transaction one attempts to acquire a row-level lock on the specified row, but it cannot: transaction two already holds such a lock. So it waits for transaction two to complete. Thus, transaction one is blocked on transaction two, and transaction two is blocked on transaction one: a deadlock condition. PostgreSQL will detect this situation and abort one of the transactions.

The best defense against deadlocks is generally to avoid them by being certain that all applications using a database acquire locks on multiple objects in a consistent order. In the example above, if both transactions had updated the rows in the same order, no deadlock would have occurred. One should also ensure that the first lock acquired on an object in a transaction is the most restrictive mode that will be needed for that object. If it is not feasible to verify this in advance, then deadlocks can be handled on-the-fly by retrying transactions that abort due to deadlocks.

So long as no deadlock situation is detected, a transaction seeking either a table-level or row-level lock will wait indefinitely for conflicting locks to be released. This means it is a bad idea for applications to hold transactions open for long periods of time (e.g., while waiting for user input).

### 13.3.5. Advisory Locks

PostgreSQL provides a means for creating locks that have application-defined meanings. These are called advisory locks, because the system does not enforce their use — it is up to the application to use them correctly. Advisory locks can be useful for locking strategies that are an awkward fit for the MVCC model. For example, a common use of advisory locks is to emulate pessimistic locking strategies typical of so-called “flat file” data management systems. While a flag stored in a table could be used for the same purpose, advisory locks are faster, avoid table bloat, and are automatically cleaned up by the server at the end of the session.

There are two ways to acquire an advisory lock in PostgreSQL: at session level or at transaction level. Once acquired at session level, an advisory lock is held until explicitly released or the session ends. Unlike standard lock requests, session-level advisory lock requests do not honor transaction semantics: a lock acquired during a transaction that is later rolled back will still be held following the rollback, and likewise an unlock is effective even if the calling transaction fails later. A lock can be acquired multiple times by its owning process; for each completed lock request there must be a corresponding unlock request before the lock is actually released. Transaction-level lock requests, on the other hand, behave more like regular lock requests: they are automatically released at the end of the transaction, and there is no explicit unlock operation. This behavior is often more convenient than the session-level behavior for short-term usage of an advisory lock. Session-level and transaction-level lock requests for the same advisory lock identifier will block each other in the expected way. If a session already holds a given advisory lock, additional requests by it will always succeed, even if other sessions are awaiting the lock; this statement is true regardless of whether the existing lock hold and new request are at session level or transaction level.
Chapter 13. Concurrency Control

Like all locks in PostgreSQL, a complete list of advisory locks currently held by any session can be found in the `pg_locks` system view.

Both advisory locks and regular locks are stored in a shared memory pool whose size is defined by the configuration variables `max_locks_per_transaction` and `max_connections`. Care must be taken not to exhaust this memory or the server will be unable to grant any locks at all. This imposes an upper limit on the number of advisory locks grantable by the server, typically in the tens to hundreds of thousands depending on how the server is configured.

In certain cases using advisory locking methods, especially in queries involving explicit ordering and `LIMIT` clauses, care must be taken to control the locks acquired because of the order in which SQL expressions are evaluated. For example:

```sql
SELECT pg_advisory_lock(id) FROM foo WHERE id = 12345; -- ok
SELECT pg_advisory_lock(id) FROM foo WHERE id > 12345 LIMIT 100; -- danger!
SELECT pg_advisory_lock(q.id) FROM (
    SELECT id FROM foo WHERE id > 12345 LIMIT 100
) q; -- ok
```

In the above queries, the second form is dangerous because the `LIMIT` is not guaranteed to be applied before the locking function is executed. This might cause some locks to be acquired that the application was not expecting, and hence would fail to release (until it ends the session). From the point of view of the application, such locks would be dangling, although still viewable in `pg_locks`.

The functions provided to manipulate advisory locks are described in Section 9.26.10.

### 13.4. Data Consistency Checks at the Application Level

It is very difficult to enforce business rules regarding data integrity using Read Committed transactions because the view of the data is shifting with each statement, and even a single statement may not restrict itself to the statement’s snapshot if a write conflict occurs.

While a Repeatable Read transaction has a stable view of the data throughout its execution, there is a subtle issue with using MVCC snapshots for data consistency checks, involving something known as read/write conflicts. If one transaction writes data and a concurrent transaction attempts to read the same data (whether before or after the write), it cannot see the work of the other transaction. The reader then appears to have executed first regardless of which started first or which committed first. If that is as far as it goes, there is no problem, but if the reader also writes data which is read by a concurrent transaction there is now a transaction which appears to have run before either of the previously mentioned transactions. If the transaction which appears to have executed last actually commits first, it is very easy for a cycle to appear in a graph of the order of execution of the transactions. When such a cycle appears, integrity checks will not work correctly without some help.

As mentioned in Section 13.2.3, Serializable transactions are just Repeatable Read transactions which add nonblocking monitoring for dangerous patterns of read/write conflicts. When a pattern is detected which could cause a cycle in the apparent order of execution, one of the transactions involved is rolled back to break the cycle.
13.4.1. Enforcing Consistency With Serializable Transactions

If the Serializable transaction isolation level is used for all writes and for all reads which need a consistent view of the data, no other effort is required to ensure consistency. Software from other environments which is written to use serializable transactions to ensure consistency should “just work” in this regard in PostgreSQL.

When using this technique, it will avoid creating an unnecessary burden for application programmers if the application software goes through a framework which automatically retries transactions which are rolled back with a serialization failure. It may be a good idea to set `default_transaction_isolation` to `serializable`. It would also be wise to take some action to ensure that no other transaction isolation level is used, either inadvertently or to subvert integrity checks, through checks of the transaction isolation level in triggers.

See Section 13.2.3 for performance suggestions.

### Warning

This level of integrity protection using Serializable transactions does not yet extend to hot standby mode (Section 26.5). Because of that, those using hot standby may want to use Repeatable Read and explicit locking on the master.

13.4.2. Enforcing Consistency With Explicit Blocking Locks

When non-serializable writes are possible, to ensure the current validity of a row and protect it against concurrent updates one must use `SELECT FOR UPDATE`, `SELECT FOR SHARE`, or an appropriate `LOCK TABLE` statement. (`SELECT FOR UPDATE` and `SELECT FOR SHARE` lock just the returned rows against concurrent updates, while `LOCK TABLE` locks the whole table.) This should be taken into account when porting applications to PostgreSQL from other environments.

Also of note to those converting from other environments is the fact that `SELECT FOR UPDATE` does not ensure that a concurrent transaction will not update or delete a selected row. To do that in PostgreSQL you must actually update the row, even if no values need to be changed. `SELECT FOR UPDATE` temporarily blocks other transactions from acquiring the same lock or executing an `UPDATE` or `DELETE` which would affect the locked row, but once the transaction holding this lock commits or rolls back, a blocked transaction will proceed with the conflicting operation unless an actual `UPDATE` of the row was performed while the lock was held.

Global validity checks require extra thought under non-serializable MVCC. For example, a banking application might wish to check that the sum of all credits in one table equals the sum of debits in another table, when both tables are being actively updated. Comparing the results of two successive `SELECT sum( ... )` commands will not work reliably in Read Committed mode, since the second query will likely include the results of transactions not counted by the first. Doing the two sums in a single repeatable read transaction will give an accurate picture of only the effects of transactions that committed before the repeatable read transaction started — but one might legitimately wonder whether the answer is still relevant by the time it is delivered. If the repeatable read transaction itself applied some changes before trying to make the consistency check, the usefulness of the check becomes even more debatable, since now it includes some but not all post-transaction-start changes. In such cases a careful person might wish to lock all tables needed for the check, in order to get an indisputable picture of current reality. A `SHARE` mode (or higher) lock guarantees that there are no uncommitted changes in the locked table, other than those of the current transaction.
Chapter 13. Concurrency Control

Note also that if one is relying on explicit locking to prevent concurrent changes, one should either use Read Committed mode, or in Repeatable Read mode be careful to obtain locks before performing queries. A lock obtained by a repeatable read transaction guarantees that no other transactions modifying the table are still running, but if the snapshot seen by the transaction predates obtaining the lock, it might predate some now-committed changes in the table. A repeatable read transaction’s snapshot is actually frozen at the start of its first query or data-modification command (SELECT, INSERT, UPDATE, or DELETE), so it is possible to obtain locks explicitly before the snapshot is frozen.

13.5. Caveats

Some DDL commands, currently only TRUNCATE and the table-rewriting forms of ALTER TABLE, are not MVCC-safe. This means that after the truncation or rewrite commits, the table will appear empty to concurrent transactions, if they are using a snapshot taken before the DDL command committed. This will only be an issue for a transaction that did not access the table in question before the DDL command started — any transaction that has done so would hold at least an ACCESS SHARE table lock, which would block the DDL command until that transaction completes. So these commands will not cause any apparent inconsistency in the table contents for successive queries on the target table, but they could cause visible inconsistency between the contents of the target table and other tables in the database.

Support for the Serializable transaction isolation level has not yet been added to Hot Standby replication targets (described in Section 26.5). The strictest isolation level currently supported in hot standby mode is Repeatable Read. While performing all permanent database writes within Serializable transactions on the master will ensure that all standbys will eventually reach a consistent state, a Repeatable Read transaction run on the standby can sometimes see a transient state that is inconsistent with any serial execution of the transactions on the master.

13.6. Locking and Indexes

Though PostgreSQL provides nonblocking read/write access to table data, nonblocking read/write access is not currently offered for every index access method implemented in PostgreSQL. The various index types are handled as follows:

B-tree, GiST and SP-GiST indexes

Short-term share/exclusive page-level locks are used for read/write access. Locks are released immediately after each index row is fetched or inserted. These index types provide the highest concurrency without deadlock conditions.

Hash indexes

Share/exclusive hash-bucket-level locks are used for read/write access. Locks are released after the whole bucket is processed. Bucket-level locks provide better concurrency than index-level ones, but deadlock is possible since the locks are held longer than one index operation.

GIN indexes

Short-term share/exclusive page-level locks are used for read/write access. Locks are released immediately after each index row is fetched or inserted. But note that insertion of a GIN-indexed value usually produces several index key insertions per row, so GIN might do substantial work for a single value’s insertion.
Currently, B-tree indexes offer the best performance for concurrent applications; since they also have more features than hash indexes, they are the recommended index type for concurrent applications that need to index scalar data. When dealing with non-scalar data, B-trees are not useful, and GiST, SP-GiST or GIN indexes should be used instead.
Chapter 14. Performance Tips

Query performance can be affected by many things. Some of these can be controlled by the user, while others are fundamental to the underlying design of the system. This chapter provides some hints about understanding and tuning PostgreSQL performance.

14.1. Using EXPLAIN

PostgreSQL devises a query plan for each query it receives. Choosing the right plan to match the query structure and the properties of the data is absolutely critical for good performance, so the system includes a complex planner that tries to choose good plans. You can use the EXPLAIN command to see what query plan the planner creates for any query. Plan-reading is an art that requires some experience to master, but this section attempts to cover the basics.

Examples in this section are drawn from the regression test database after doing a VACUUM ANALYZE, using 9.3 development sources. You should be able to get similar results if you try the examples yourself, but your estimated costs and row counts might vary slightly because ANALYZE’s statistics are random samples rather than exact, and because costs are inherently somewhat platform-dependent.

The examples use EXPLAIN’s default “text” output format, which is compact and convenient for humans to read. If you want to feed EXPLAIN’s output to a program for further analysis, you should use one of its machine-readable output formats (XML, JSON, or YAML) instead.

14.1.1. EXPLAIN Basics

The structure of a query plan is a tree of plan nodes. Nodes at the bottom level of the tree are scan nodes: they return raw rows from a table. There are different types of scan nodes for different table access methods: sequential scans, index scans, and bitmap index scans. There are also non-table row sources, such as VALUES clauses and set-returning functions in FROM, which have their own scan node types. If the query requires joining, aggregation, sorting, or other operations on the raw rows, then there will be additional nodes above the scan nodes to perform these operations. Again, there is usually more than one possible way to do these operations, so different node types can appear here too. The output of EXPLAIN has one line for each node in the plan tree, showing the basic node type plus the cost estimates that the planner made for the execution of that plan node. Additional lines might appear, indented from the node’s summary line, to show additional properties of the node. The very first line (the summary line for the topmost node) has the estimated total execution cost for the plan; it is this number that the planner seeks to minimize.

Here is a trivial example, just to show what the output looks like:

```
EXPLAIN SELECT * FROM tenk1;
```

```
QUERY PLAN
-------------------------------------------------------------
Seq Scan on tenk1  (cost=0.00..458.00 rows=10000 width=244)
```

Since this query has no WHERE clause, it must scan all the rows of the table, so the planner has chosen to use a simple sequential scan plan. The numbers that are quoted in parentheses are (left to right):
Chapter 14. Performance Tips

- Estimated start-up cost. This is the time expended before the output phase can begin, e.g., time to do the sorting in a sort node.

- Estimated total cost. This is stated on the assumption that the plan node is run to completion, i.e., all available rows are retrieved. In practice a node’s parent node might stop short of reading all available rows (see the LIMIT example below).

- Estimated number of rows output by this plan node. Again, the node is assumed to be run to completion.

- Estimated average width of rows output by this plan node (in bytes).

The costs are measured in arbitrary units determined by the planner’s cost parameters (see Section 19.7.2). Traditional practice is to measure the costs in units of disk page fetches; that is, seq_page_cost is conventionally set to 1.0 and the other cost parameters are set relative to that. The examples in this section are run with the default cost parameters.

It’s important to understand that the cost of an upper-level node includes the cost of all its child nodes. It’s also important to realize that the cost only reflects things that the planner cares about. In particular, the cost does not consider the time spent transmitting result rows to the client, which could be an important factor in the real elapsed time; but the planner ignores it because it cannot change it by altering the plan. (Every correct plan will output the same row set, we trust.)

The rows value is a little tricky because it is not the number of rows processed or scanned by the plan node, but rather the number emitted by the node. This is often less than the number scanned, as a result of filtering by any WHERE-clause conditions that are being applied at the node. Ideally the top-level rows estimate will approximate the number of rows actually returned, updated, or deleted by the query.

Returning to our example:

```
EXPLAIN SELECT * FROM tenk1;
```

```
QUERY PLAN

Seq Scan on tenk1  (cost=0.00..458.00 rows=10000 width=244)

```

These numbers are derived very straightforwardly. If you do:

```
SELECT relpages, reltuples FROM pg_class WHERE relname = 'tenk1';
```

you will find that tenk1 has 358 disk pages and 10000 rows. The estimated cost is computed as (disk pages read * seq_page_cost) + (rows scanned * cpu_tuple_cost). By default, seq_page_cost is 1.0 and cpu_tuple_cost is 0.01, so the estimated cost is (358 * 1.0) + (10000 * 0.01) = 458.

Now let’s modify the query to add a WHERE condition:

```
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 7000;
```

```
QUERY PLAN

Seq Scan on tenk1  (cost=0.00..483.00 rows=7001 width=244)
Filter: (unique1 < 7000)

```

Notice that the EXPLAIN output shows the WHERE clause being applied as a “filter” condition attached to the Seq Scan plan node. This means that the plan node checks the condition for each row it scans,
and outputs only the ones that pass the condition. The estimate of output rows has been reduced because of the `WHERE` clause. However, the scan will still have to visit all 10000 rows, so the cost hasn’t decreased; in fact it has gone up a bit (by 10000 * cpu_operator_cost, to be exact) to reflect the extra CPU time spent checking the `WHERE` condition.

The actual number of rows this query would select is 7000, but the `ROWS` estimate is only approximate. If you try to duplicate this experiment, you will probably get a slightly different estimate; moreover, it can change after each `ANALYZE` command, because the statistics produced by `ANALYZE` are taken from a randomized sample of the table.

Now, let’s make the condition more restrictive:

```sql
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100;
```

### QUERY PLAN
```
-----------------------------------------------------------------------------
Bitmap Heap Scan on tenk1 (cost=5.07..229.20 rows=101 width=244)
  Recheck Cond: (unique1 < 100)
--> Bitmap Index Scan on tenk1_unique1 (cost=0.00..5.04 rows=101 width=0)
    Index Cond: (unique1 < 100)
-----------------------------------------------------------------------------
```

Here the planner has decided to use a two-step plan: the child plan node visits an index to find the locations of rows matching the index condition, and then the upper plan node actually fetches those rows from the table itself. Fetching rows separately is much more expensive than reading them sequentially, but because not all the pages of the table have to be visited, this is still cheaper than a sequential scan. (The reason for using two plan levels is that the upper plan node sorts the row locations identified by the index into physical order before reading them, to minimize the cost of separate fetches. The “bitmap” mentioned in the node names is the mechanism that does the sorting.)

Now let’s add another condition to the `WHERE` clause:

```sql
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100 AND stringu1 = 'xxx';
```

### QUERY PLAN
```
-----------------------------------------------------------------------------
Bitmap Heap Scan on tenk1 (cost=5.04..229.43 rows=1 width=244)
  Recheck Cond: (unique1 < 100)
  Filter: (stringu1 = 'xxx'::name)
--> Bitmap Index Scan on tenk1_unique1 (cost=0.00..5.04 rows=101 width=0)
    Index Cond: (unique1 < 100)
-----------------------------------------------------------------------------
```

The added condition `stringu1 = 'xxx'` reduces the output row count estimate, but not the cost because we still have to visit the same set of rows. Notice that the `stringu1` clause cannot be applied as an index condition, since this index is only on the `unique1` column. Instead it is applied as a filter on the rows retrieved by the index. Thus the cost has actually gone up slightly to reflect this extra checking.

In some cases the planner will prefer a “simple” index scan plan:

```sql
EXPLAIN SELECT * FROM tenk1 WHERE unique1 = 42;
```

### QUERY PLAN
```
-----------------------------------------------------------------------------
Index Scan using tenk1_unique1 on tenk1 (cost=0.29..8.30 rows=1 width=244)
  Index Cond: (unique1 = 42)
-----------------------------------------------------------------------------
```

In this type of plan the table rows are fetched in index order, which makes them even more expensive to read, but there are so few that the extra cost of sorting the row locations is not worth it. You’ll most
often see this plan type for queries that fetch just a single row. It’s also often used for queries that have an ORDER BY condition that matches the index order, because then no extra sorting step is needed to satisfy the ORDER BY.

If there are separate indexes on several of the columns referenced in WHERE, the planner might choose to use an AND or OR combination of the indexes:

EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000;

QUERY PLAN

- Bitmap Heap Scan on tenk1 (cost=25.08..60.21 rows=10 width=244)
  - Recheck Cond: ((unique1 < 100) AND (unique2 > 9000))
    - BitmapAnd (cost=25.08..25.08 rows=10 width=0)
      - Bitmap Index Scan on tenk1_unique1 (cost=0.00..5.04 rows=101 width=0)
        - Index Cond: (unique1 < 100)
      - Bitmap Index Scan on tenk1_unique2 (cost=0.00..19.78 rows=999 width=0)
        - Index Cond: (unique2 > 9000)

But this requires visiting both indexes, so it’s not necessarily a win compared to using just one index and treating the other condition as a filter. If you vary the ranges involved you’ll see the plan change accordingly.

Here is an example showing the effects of LIMIT:

EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000 LIMIT 2;

QUERY PLAN

- Limit (cost=0.29..14.48 rows=2 width=244)
  - Index Scan using tenk2_unique2 on tenk2 (cost=0.29..71.27 rows=1 width=244)
    - Index Cond: (unique2 > 9000)
    - Filter: (unique1 < 100)

This is the same query as above, but we added a LIMIT so that not all the rows need be retrieved, and the planner changed its mind about what to do. Notice that the total cost and row count of the Index Scan node are shown as if it were run to completion. However, the Limit node is expected to stop after retrieving only a fifth of those rows, so its total cost is only a fifth as much, and that’s the actual estimated cost of the query. This plan is preferred over adding a Limit node to the previous plan because the Limit could not avoid paying the startup cost of the bitmap scan, so the total cost would be something over 25 units with that approach.

Let’s try joining two tables, using the columns we have been discussing:

EXPLAIN SELECT * FROM tenk1 t1, tenk2 t2 WHERE t1.unique1 < 10 AND t1.unique2 = t2.unique2;

QUERY PLAN

- Nested Loop (cost=4.65..118.62 rows=10 width=488)
  - Bitmap Heap Scan on tenk1 t1 (cost=4.36..39.47 rows=10 width=244)
    - Recheck Cond: (unique1 < 10)
      - Bitmap Index Scan on tenk1_unique1 (cost=0.00..4.36 rows=10 width=0)
        - Index Cond: (unique1 < 10)
      - Index Scan using tenk2_unique2 on tenk2 t2 (cost=0.29..7.91 rows=1 width=244)
In this plan, we have a nested-loop join node with two table scans as inputs, or children. The indentation of the node summary lines reflects the plan tree structure. The join’s first, or “outer”, child is a bitmap scan similar to those we saw before. Its cost and row count are the same as we’d get from

\[
\text{SELECT ... WHERE unique1 < 10 because we are applying the WHERE clause unique1 < 10 at that node.}
\]

The `t1.unique2 = t2.unique2` clause is not relevant yet, so it doesn’t affect the row count of the outer scan. The nested-loop join node will run its second, or “inner” child once for each row obtained from the outer child. Column values from the current outer row can be plugged into the inner scan; here, the `t1.unique2` value from the outer row is available, so we get a plan and costs similar to what we saw above for a simple `SELECT ... WHERE t2.unique2 = constant` case.

(The estimated cost is actually a bit lower than what was seen above, as a result of caching that’s expected to occur during the repeated index scans on `t2`.) The costs of the loop node are then set on the basis of the cost of the outer scan, plus one repetition of the inner scan for each outer row (10 * 7.91, here), plus a little CPU time for join processing.

In this example the join’s output row count is the same as the product of the two scans’ row counts, but that’s not true in all cases because there can be additional `WHERE` clauses that mention both tables and so can only be applied at the join point, not to either input scan. Here’s an example:

```
EXPLAIN SELECT *
FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 10 AND t2.unique2 < 10 AND t1.hundred < t2.hundred;
```

```
QUERY PLAN

---------------------------------------------------------------------------------------------
| Nested Loop (cost=4.65..49.46 rows=33 width=488) |
| Join Filter: (t1.hundred < t2.hundred) |
| -> Bitmap Heap Scan on tenk1 t1 (cost=4.36..39.47 rows=10 width=244) |
|    Recheck Cond: (unique1 < 10) |
|    -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..4.36 rows=10 width=0) |
|    Index Cond: (unique1 < 10) |
| -> Materialize (cost=0.29..8.51 rows=10 width=244) |
|    -> Index Scan using tenk2_unique2 on tenk2 t2 (cost=0.29..8.46 rows=10 width=0) |
|    Index Cond: (unique2 < 10) |
| The condition t1.hundred < t2.hundred can’t be tested in the tenk2_unique2 index, so it’s applied at the join node. This reduces the estimated output row count of the join node, but does not change either input scan. |
| Notice that here the planner has chosen to “materialize” the inner relation of the join, by putting a Materialize plan node atop it. This means that the `t2` index scan will be done just once, even though the nested-loop join node needs to read that data ten times, once for each row from the outer relation. The Materialize node saves the data in memory as it’s read, and then returns the data from memory on each subsequent pass. |
| When dealing with outer joins, you might see join plan nodes with both “Join Filter” and plain “Filter” conditions attached. Join Filter conditions come from the outer join’s `ON` clause, so a row that fails the Join Filter condition could still get emitted as a null-extended row. But a plain Filter condition is applied after the outer-join rules and so acts to remove rows unconditionally. In an inner join there is no semantic difference between these types of filters. |
| If we change the query’s selectivity a bit, we might get a very different join plan: |
```

EXPLAIN SELECT *
```
FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2;

QUERY PLAN
------------------------------------------------------------------------------------------
Hash Join (cost=230.47..713.98 rows=101 width=488)
  Hash Cond: (t2.unique2 = t1.unique2)
    ->  Seq Scan on tenk2 t2 (cost=0.00..445.00 rows=10000 width=244)
    ->  Hash (cost=229.20..229.20 rows=101 width=244)
       -> Bitmap Heap Scan on tenk1 t1 (cost=5.07..229.20 rows=101 width=244)
          Recheck Cond: (unique1 < 100)
          -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..5.04 rows=101 width=0)
             Index Cond: (unique1 < 100)

Here, the planner has chosen to use a hash join, in which rows of one table are entered into an in-memory hash table, after which the other table is scanned and the hash table is probed for matches to each row. Again note how the indentation reflects the plan structure: the bitmap scan on tenk1 is the input to the Hash node, which constructs the hash table. That’s then returned to the Hash Join node, which reads rows from its outer child plan and searches the hash table for each one.

Another possible type of join is a merge join, illustrated here:

EXPLAIN SELECT *
FROM tenk1 t1, onek t2
WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2;

QUERY PLAN
------------------------------------------------------------------------------------------
Merge Join (cost=198.11..268.19 rows=10 width=488)
  Merge Cond: (t1.unique2 = t2.unique2)
    ->  Index Scan using tenk1_unique2 on tenk1 t1 (cost=0.29..656.28 rows=101 width=244)
        Filter: (unique1 < 100)
    ->  Sort (cost=197.83..200.33 rows=1000 width=244)
        Sort Key: t2.unique2
        ->  Seq Scan on onek t2 (cost=0.00..148.00 rows=1000 width=244)

Merge join requires its input data to be sorted on the join keys. In this plan the tenk1 data is sorted by using an index scan to visit the rows in the correct order, but a sequential scan and sort is preferred for onek, because there are many more rows to be visited in that table. (Sequential-scan-and-sort frequently beats an index scan for sorting many rows, because of the nonsequential disk access required by the index scan.)

One way to look at variant plans is to force the planner to disregard whatever strategy it thought was the cheapest, using the enable/disable flags described in Section 19.7.1. (This is a crude tool, but useful. See also Section 14.3.) For example, if we’re unconvinced that sequential-scan-and-sort is the best way to deal with table onek in the previous example, we could try

SET enable_sort = off;

EXPLAIN SELECT *
FROM tenk1 t1, onek t2
WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2;

QUERY PLAN
which shows that the planner thinks that sorting onek by index-scanning is about 12% more expensive than sequential-scan-and-sort. Of course, the next question is whether it’s right about that. We can investigate that using EXPLAIN ANALYZE, as discussed below.

### 14.1.2. EXPLAIN ANALYZE

It is possible to check the accuracy of the planner’s estimates by using EXPLAIN’s ANALYZE option. With this option, EXPLAIN actually executes the query, and then displays the true row counts and true run time accumulated within each plan node, along with the same estimates that a plain EXPLAIN shows. For example, we might get a result like this:

```
EXPLAIN ANALYZE SELECT * FROM tenk1 t1, tenk2 t2 WHERE t1.unique1 < 10 AND t1.unique2 = t2.unique2;
```

```sql
QUERY PLAN
---------------------------------------------------------------------------------------------------------------------------------
Nested Loop (cost=4.65..118.62 rows=10 width=488) (actual time=0.128..0.377 rows=10 loops=1)
  -> Bitmap Heap Scan on tenk1 t1 (cost=4.36..39.47 rows=10 width=244) (actual time=0.057..0.121 rows=10 loops=1)
    Recheck Cond: (unique1 < 10)
    -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..4.36 rows=10 width=0) (actual time=0.024..0.024 rows=10 loops=1)
      Index Cond: (unique1 < 10)
  -> Index Scan using tenk2_unique2 on tenk2 t2 (cost=0.29..224.79 rows=1000 width=244) (actual time=0.021..0.022 rows=1 loops=10)
    Index Cond: (unique2 = t1.unique2)
Planning time: 0.181 ms
Execution time: 0.501 ms
```

Note that the “actual time” values are in milliseconds of real time, whereas the cost estimates are expressed in arbitrary units; so they are unlikely to match up. The thing that’s usually most important to look for is whether the estimated row counts are reasonably close to reality. In this example the estimates were all dead-on, but that’s quite unusual in practice.

In some query plans, it is possible for a subplan node to be executed more than once. For example, the inner index scan will be executed once per outer row in the above nested-loop plan. In such cases, the loops value reports the total number of executions of the node, and the actual time and rows values shown are averages per-execution. This is done to make the numbers comparable with the way that the cost estimates are shown. Multiply by the loops value to get the total time actually spent in the node. In the above example, we spent a total of 0.220 milliseconds executing the index scans on tenk2.

In some cases EXPLAIN ANALYZE shows additional execution statistics beyond the plan node execution times and row counts. For example, Sort and Hash nodes provide extra information:

```
EXPLAIN ANALYZE SELECT * FROM tenk1 t1, tenk2 t2 WHERE t1.unique1 < 100 AND t1.unique2 = t2.unique2 ORDER BY t1.fivethous;
```

```sql
QUERY PLAN
---------------------------------------------------------------------------------------------------------------------------------
Nested Loop (cost=4.65..118.62 rows=10 width=488) (actual time=0.128..0.377 rows=10 loops=1)
  -> Bitmap Heap Scan on tenk1 t1 (cost=4.36..39.47 rows=10 width=244) (actual time=0.057..0.121 rows=10 loops=1)
    Recheck Cond: (unique1 < 10)
    -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..4.36 rows=10 width=0) (actual time=0.024..0.024 rows=10 loops=1)
      Index Cond: (unique1 < 10)
  -> Index Scan using tenk2_unique2 on tenk2 t2 (cost=0.29..224.79 rows=1000 width=244) (actual time=0.021..0.022 rows=1 loops=10)
    Index Cond: (unique2 = t1.unique2)
Planning time: 0.181 ms
Execution time: 0.501 ms
```
Chapter 14. Performance Tips

Sort (cost=717.34..717.59 rows=101 width=488) (actual time=7.761..7.774 rows=100 loops=1)
Sort Key: t1.fivethous
Sort Method: quicksort Memory: 77kB
-> Hash Join (cost=230.47..713.98 rows=101 width=488) (actual time=0.711..7.427 rows=100 loops=1)
  Hash Cond: (t2.unique2 = t1.unique2)
  -> Seq Scan on tenk2 t2 (cost=0.00..445.00 rows=10000 width=244) (actual time=0.007..2.583 rows=10000 loops=1)
  Hash (cost=229.20..229.20 rows=101 width=244) (actual time=0.659..0.659 rows=100 loops=1)
    Buckets: 1024 Batches: 1 Memory Usage: 28kB
    -> Bitmap Heap Scan on tenk1 t1 (cost=5.07..229.20 rows=101 width=244)
       Recheck Cond: (unique1 < 100)
       -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..5.04 rows=101 width=0)
          Index Cond: (unique1 < 100)
Planning time: 0.194 ms
Execution time: 8.008 ms

The Sort node shows the sort method used (in particular, whether the sort was in-memory or on-disk) and the amount of memory or disk space needed. The Hash node shows the number of hash buckets and batches as well as the peak amount of memory used for the hash table. (If the number of batches exceeds one, there will also be disk space usage involved, but that is not shown.)

Another type of extra information is the number of rows removed by a filter condition:

EXPLAIN ANALYZE SELECT * FROM tenk1 WHERE ten < 7;

QUERY PLAN

Seq Scan on tenk1 (cost=0.00..483.00 rows=7000 width=244) (actual time=0.016..5.107 rows=7000 loops=1)
  Filter: (ten < 7)
  Rows Removed by Filter: 3000
Planning time: 0.083 ms
Execution time: 5.905 ms

These counts can be particularly valuable for filter conditions applied at join nodes. The “Rows Removed” line only appears when at least one scanned row, or potential join pair in the case of a join node, is rejected by the filter condition.

A case similar to filter conditions occurs with “lossy” index scans. For example, consider this search for polygons containing a specific point:

EXPLAIN ANALYZE SELECT * FROM polygon_tbl WHERE f1 @> polygon '(0.5,2.0)';

QUERY PLAN

Seq Scan on polygon_tbl (cost=0.00..1.05 rows=1 width=32) (actual time=0.044..0.044 rows=0 loops=1)
  Filter: (f1 @> '((0.5,2))':polygon)
  Rows Removed by Filter: 4
Planning time: 0.040 ms
Execution time: 0.083 ms

The planner thinks (quite correctly) that this sample table is too small to bother with an index scan, so we have a plain sequential scan in which all the rows got rejected by the filter condition. But if we force an index scan to be used, we see:

SET enable_seqscan TO off;

EXPLAIN ANALYZE SELECT * FROM polygon_tbl WHERE f1 @> polygon '(0.5,2.0)';
Here we can see that the index returned one candidate row, which was then rejected by a recheck of the index condition. This happens because a GiST index is “lossy” for polygon containment tests: it actually returns the rows with polygons that overlap the target, and then we have to do the exact containment test on those rows.

EXPLAIN has a BUFFERS option that can be used with ANALYZE to get even more run time statistics:

EXPLAIN (ANALYZE, BUFFERS) SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000;

The numbers provided by BUFFERS help to identify which parts of the query are the most I/O-intensive.

Keep in mind that because EXPLAIN ANALYZE actually runs the query, any side-effects will happen as usual, even though whatever results the query might output are discarded in favor of printing the EXPLAIN data. If you want to analyze a data-modifying query without changing your tables, you can roll the command back afterwards, for example:

BEGIN;
EXPLAIN ANALYZE UPDATE tenk1 SET hundred = hundred + 1 WHERE unique1 < 100;
ROLLBACK;
As seen in this example, when the query is an **INSERT**, **UPDATE**, or **DELETE** command, the actual work of applying the table changes is done by a top-level Insert, Update, or Delete plan node. The plan nodes underneath this node perform the work of locating the old rows and/or computing the new data. So above, we see the same sort of bitmap table scan we’ve seen already, and its output is fed to an Update node that stores the updated rows. It’s worth noting that although the data-modifying node can take a considerable amount of run time (here, it’s consuming the lion’s share of the time), the planner does not currently add anything to the cost estimates to account for that work. That’s because the work to be done is the same for every correct query plan, so it doesn’t affect planning decisions.

When an **UPDATE** or **DELETE** command affects an inheritance hierarchy, the output might look like this:

```
EXPLAIN UPDATE parent SET f2 = f2 + 1 WHERE f1 = 101;
```

```
QUERY PLAN
-----------------------------------------------------------------------------------
Update on parent (cost=0.00..24.53 rows=4 width=14)
  Update on parent
  Update on child1
  Update on child2
  Update on child3
    -> Seq Scan on parent (cost=0.00..0.00 rows=1 width=14)
      Filter: (f1 = 101)
    -> Index Scan using child1_f1_key on child1 (cost=0.15..8.17 rows=1 width=14)
      Index Cond: (f1 = 101)
    -> Index Scan using child2_f1_key on child2 (cost=0.15..8.17 rows=1 width=14)
      Index Cond: (f1 = 101)
    -> Index Scan using child3_f1_key on child3 (cost=0.15..8.17 rows=1 width=14)
      Index Cond: (f1 = 101)
```

In this example the Update node needs to consider three child tables as well as the originally-mentioned parent table. So there are four input scanning subplans, one per table. For clarity, the Update node is annotated to show the specific target tables that will be updated, in the same order as the corresponding subplans. (These annotations are new as of PostgreSQL 9.5; in prior versions the reader had to intuit the target tables by inspecting the subplans.)

The **Planning time** shown by **EXPLAIN ANALYZE** is the time it took to generate the query plan from the parsed query and optimize it. It does not include parsing or rewriting.

The **Execution time** shown by **EXPLAIN ANALYZE** includes executor start-up and shut-down time, as well as the time to run any triggers that are fired, but it does not include parsing, rewriting, or planning time. Time spent executing **BEFORE** triggers, if any, is included in the time for the related Insert, Update, or Delete node; but time spent executing **AFTER** triggers is not counted there because **AFTER** triggers are fired after completion of the whole plan. The total time spent in each trigger (either **BEFORE** or **AFTER**) is also shown separately. Note that deferred constraint triggers will not be executed until end of transaction and are thus not considered at all by **EXPLAIN ANALYZE**.

### 14.1.3. Caveats

There are two significant ways in which run times measured by **EXPLAIN ANALYZE** can deviate from normal execution of the same query. First, since no output rows are delivered to the client, network transmission costs and I/O conversion costs are not included. Second, the measurement overhead added by **EXPLAIN ANALYZE** can be significant, especially on machines with slow `gettimeofday()` operating-system calls. You can use the `pg_test_timing` tool to measure the overhead of timing on your system.
EXPLAIN results should not be extrapolated to situations much different from the one you are actually testing; for example, results on a toy-sized table cannot be assumed to apply to large tables. The planner’s cost estimates are not linear and so it might choose a different plan for a larger or smaller table. An extreme example is that on a table that only occupies one disk page, you’ll nearly always get a sequential scan plan whether indexes are available or not. The planner realizes that it’s going to take one disk page read to process the table in any case, so there’s no value in expending additional page reads to look at an index. (We saw this happening in the polygon_tbl example above.)

There are cases in which the actual and estimated values won’t match up well, but nothing is really wrong. One such case occurs when plan node execution is stopped short by a LIMIT or similar effect. For example, in the LIMIT query we used before,

```
EXPLAIN ANALYZE SELECT * FROM tenk1 WHERE unique1 < 100 AND unique2 > 9000 LIMIT 2;
```

the estimated cost and row count for the Index Scan node are shown as though it were run to completion. But in reality the Limit node stopped requesting rows after it got two, so the actual row count is only 2 and the run time is less than the cost estimate would suggest. This is not an estimation error, only a discrepancy in the way the estimates and true values are displayed.

Merge joins also have measurement artifacts that can confuse the unwary. A merge join will stop reading one input if it’s exhausted the other input and the next key value in the one input is greater than the last key value of the other input; in such a case there can be no more matches and so no need to scan the rest of the first input. This results in not reading all of one child, with results like those mentioned for LIMIT. Also, if the outer (first) child contains rows with duplicate key values, the inner (second) child is backed up and rescanned for the portion of its rows matching that key value. EXPLAIN ANALYZE counts these repeated emissions of the same inner rows as if they were real additional rows. When there are many outer duplicates, the reported actual row count for the inner child plan node can be significantly larger than the number of rows that are actually in the inner relation.

BitmapAnd and BitmapOr nodes always report their actual row counts as zero, due to implementation limitations.

### 14.2. Statistics Used by the Planner

As we saw in the previous section, the query planner needs to estimate the number of rows retrieved by a query in order to make good choices of query plans. This section provides a quick look at the statistics that the system uses for these estimates.

One component of the statistics is the total number of entries in each table and index, as well as the number of disk blocks occupied by each table and index. This information is kept in the table `pg_class`, in the columns `reltuples` and `replpages`. We can look at it with queries similar to this one:
SELECT relname, relkind, reltuples, relpages
FROM pg_class
WHERE relname LIKE 'tenk1%';

<table>
<thead>
<tr>
<th>relname</th>
<th>relkind</th>
<th>reltuples</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>tenk1</td>
<td>r</td>
<td>10000</td>
<td>358</td>
</tr>
<tr>
<td>tenk1_hundred</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
<tr>
<td>tenk1_thous_tenthous</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
<tr>
<td>tenk1_unique1</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
<tr>
<td>tenk1_unique2</td>
<td>i</td>
<td>10000</td>
<td>30</td>
</tr>
</tbody>
</table>

(5 rows)

Here we can see that tenk1 contains 10000 rows, as do its indexes, but the indexes are (unsurprisingly) much smaller than the table.

For efficiency reasons, reltuples and relpages are not updated on-the-fly, and so they usually contain somewhat out-of-date values. They are updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE INDEX. A VACUUM or ANALYZE operation that does not scan the entire table (which is commonly the case) will incrementally update the reltuples count on the basis of the part of the table it did scan, resulting in an approximate value. In any case, the planner will scale the values it finds in pg_class to match the current physical table size, thus obtaining a closer approximation.

Most queries retrieve only a fraction of the rows in a table, due to WHERE clauses that restrict the rows to be examined. The planner thus needs to make an estimate of the selectivity of WHERE clauses, that is, the fraction of rows that match each condition in the WHERE clause. The information used for this task is stored in the pg_statistic system catalog. Entries in pg_statistic are updated by the ANALYZE and VACUUM ANALYZE commands, and are always approximate even when freshly updated.

Rather than look at pg_statistic directly, it's better to look at its view pg_stats when examining the statistics manually. pg_stats is designed to be more easily readable. Furthermore, pg_stats is readable by all, whereas pg_statistic is only readable by a superuser. (This prevents unprivileged users from learning something about the contents of other people’s tables from the statistics. The pg_stats view is restricted to show only rows about tables that the current user can read.) For example, we might do:

```sql
SELECT attname, inherited, n_distinct,
       array_to_string(most_common_vals, E'\n') as most_common_vals
FROM pg_stats
WHERE tablename = 'road';
```

<table>
<thead>
<tr>
<th>attname</th>
<th>inherited</th>
<th>n_distinct</th>
<th>most_common_vals</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>f</td>
<td>-0.363388</td>
<td>I- 580 Ramp+</td>
</tr>
<tr>
<td></td>
<td></td>
<td>I- 880</td>
<td>Ramp+</td>
</tr>
<tr>
<td></td>
<td></td>
<td>Sp Railroad</td>
<td>+</td>
</tr>
<tr>
<td></td>
<td></td>
<td>I- 580</td>
<td>+</td>
</tr>
<tr>
<td></td>
<td></td>
<td>I- 680</td>
<td>Ramp</td>
</tr>
<tr>
<td>name</td>
<td>t</td>
<td>-0.284859</td>
<td>I- 880 Ramp+</td>
</tr>
<tr>
<td></td>
<td></td>
<td>I- 580</td>
<td>Ramp+</td>
</tr>
<tr>
<td></td>
<td></td>
<td>I- 680</td>
<td>Ramp+</td>
</tr>
<tr>
<td></td>
<td></td>
<td>I- 580</td>
<td>+</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>State Hwy 13 Ramp</td>
</tr>
</tbody>
</table>

(2 rows)
Note that two rows are displayed for the same column, one corresponding to the complete inheritance hierarchy starting at the `road` table (`inherited=t`), and another one including only the `road` table itself (`inherited=f`). The amount of information stored in `pg_statistic` by `ANALYZE`, in particular the maximum number of entries in the `most_common_vals` and `histogram_bounds` arrays for each column, can be set on a column-by-column basis using the `ALTER TABLE SET STATISTICS` command, or globally by setting the `default_statistics_target` configuration variable. The default limit is presently 100 entries. Raising the limit might allow more accurate planner estimates to be made, particularly for columns with irregular data distributions, at the price of consuming more space in `pg_statistic` and slightly more time to compute the estimates. Conversely, a lower limit might be sufficient for columns with simple data distributions.

Further details about the planner’s use of statistics can be found in Chapter 67.

### 14.3. Controlling the Planner with Explicit `JOIN` Clauses

It is possible to control the query planner to some extent by using the explicit `JOIN` syntax. To see why this matters, we first need some background.

In a simple join query, such as:

```
SELECT * FROM a, b, c WHERE a.id = b.id AND b.ref = c.id;
```

the planner is free to join the given tables in any order. For example, it could generate a query plan that joins A to B, using the `WHERE` condition `a.id = b.id`, and then joins C to this joined table, using the other `WHERE` condition. Or it could join B to C and then join A to that result. Or it could join A to C and then join them with B — but that would be inefficient, since the full Cartesian product of A and C would have to be formed, there being no applicable condition in the `WHERE` clause to allow optimization of the join. (All joins in the PostgreSQL executor happen between two input tables, so it’s necessary to build up the result in one or another of these fashions.) The important point is that these different join possibilities give semantically equivalent results but might have hugely different execution costs. Therefore, the planner will explore all of them to try to find the most efficient query plan.

When a query only involves two or three tables, there aren’t many join orders to worry about. But the number of possible join orders grows exponentially as the number of tables expands. Beyond ten or so input tables it’s no longer practical to do an exhaustive search of all the possibilities, and even for six or seven tables planning might take an annoyingly long time. When there are too many input tables, the PostgreSQL planner will switch from exhaustive search to a genetic probabilistic search through a limited number of possibilities. (The switch-over threshold is set by the `geqo_threshold` run-time parameter.) The genetic search takes less time, but it won’t necessarily find the best possible plan.

When the query involves outer joins, the planner has less freedom than it does for plain (inner) joins. For example, consider:

```
SELECT * FROM a LEFT JOIN (b JOIN c ON (b.ref = c.id)) ON (a.id = b.id);
```

Although this query’s restrictions are superficially similar to the previous example, the semantics are different because a row must be emitted for each row of A that has no matching row in the join of B and C. Therefore the planner has no choice of join order here: it must join B to C and then join A to that result. Accordingly, this query takes less time to plan than the previous query. In other cases, the planner might be able to determine that more than one join order is safe. For example, given:

```
SELECT * FROM a LEFT JOIN b ON (a.bid = b.id) LEFT JOIN c ON (a.cid = c.id);
```
it is valid to join A to either B or C first. Currently, only **FULL JOIN** completely constrains the join order. Most practical cases involving **LEFT JOIN** or **RIGHT JOIN** can be rearranged to some extent. Explicit inner join syntax (**INNER JOIN**, **CROSS JOIN**, or unadorned **JOIN**) is semantically the same as listing the input relations in **FROM**, so it does not constrain the join order.

Even though most kinds of **JOIN** don’t completely constrain the join order, it is possible to instruct the PostgreSQL query planner to treat all **JOIN** clauses as constraining the join order anyway. For example, these three queries are logically equivalent:

```sql
SELECT * FROM a, b, c WHERE a.id = b.id AND b.ref = c.id;
SELECT * FROM a CROSS JOIN b CROSS JOIN c WHERE a.id = b.id AND b.ref = c.id;
SELECT * FROM a JOIN (b JOIN c ON (b.ref = c.id)) ON (a.id = b.id);
```

But if we tell the planner to honor the **JOIN** order, the second and third take less time to plan than the first. This effect is not worth worrying about for only three tables, but it can be a lifesaver with many tables.

To force the planner to follow the join order laid out by explicit **JOIN**s, set the **join_collapse_limit** run-time parameter to 1. (Other possible values are discussed below.) You do not need to constrain the join order completely in order to cut search time, because it’s OK to use **JOIN** operators within items of a plain **FROM** list. For example, consider:

```sql
SELECT * FROM a CROSS JOIN b, c, d, e WHERE ...;
```

With **join_collapse_limit = 1**, this forces the planner to join A to B before joining them to other tables, but doesn’t constrain its choices otherwise. In this example, the number of possible join orders is reduced by a factor of 5.

Constraining the planner’s search in this way is a useful technique both for reducing planning time and for directing the planner to a good query plan. If the planner chooses a bad join order by default, you can force it to choose a better order via **JOIN** syntax — assuming that you know of a better order, that is. Experimentation is recommended.

A closely related issue that affects planning time is collapsing of subqueries into their parent query. For example, consider:

```sql
SELECT * FROM x, y,
    (SELECT * FROM a, b, c WHERE something) AS ss
WHERE somethingelse;
```

This situation might arise from use of a view that contains a join; the view’s **SELECT** rule will be inserted in place of the view reference, yielding a query much like the above. Normally, the planner will try to collapse the subquery into the parent, yielding:

```sql
SELECT * FROM x, y, a, b, c WHERE something AND somethingelse;
```

This usually results in a better plan than planning the subquery separately. (For example, the outer **WHERE** conditions might be such that joining X to A first eliminates many rows of A, thus avoiding the need to form the full logical output of the subquery.) But at the same time, we have increased the planning time: here, we have a five-way join problem replacing two separate three-way join problems. Because of the exponential growth of the number of possibilities, this makes a big difference. The planner tries to avoid getting stuck in huge join search problems by not collapsing a subquery if more than **from_collapse_limit** **FROM** items would result in the parent query. You can trade off planning time against quality of plan by adjusting this run-time parameter up or down.
from_collapse_limit and join_collapse_limit are similarly named because they do almost the same thing: one controls when the planner will “flatten out” subqueries, and the other controls when it will flatten out explicit joins. Typically you would either set join_collapse_limit equal to from_collapse_limit (so that explicit joins and subqueries act similarly) or set join_collapse_limit to 1 (if you want to control join order with explicit joins). But you might set them differently if you are trying to fine-tune the trade-off between planning time and run time.

14.4. Populating a Database

One might need to insert a large amount of data when first populating a database. This section contains some suggestions on how to make this process as efficient as possible.

14.4.1. Disable Autocommit

When using multiple INSERTs, turn off autocommit and just do one commit at the end. (In plain SQL, this means issuing BEGIN at the start and COMMIT at the end. Some client libraries might do this behind your back, in which case you need to make sure the library does it when you want it done.) If you allow each insertion to be committed separately, PostgreSQL is doing a lot of work for each row that is added. An additional benefit of doing all insertions in one transaction is that if the insertion of one row were to fail then the insertion of all rows inserted up to that point would be rolled back, so you won’t be stuck with partially loaded data.

14.4.2. Use COPY

Use COPY to load all the rows in one command, instead of using a series of INSERT commands. The COPY command is optimized for loading large numbers of rows; it is less flexible than INSERT, but incurs significantly less overhead for large data loads. Since COPY is a single command, there is no need to disable autocommit if you use this method to populate a table.

If you cannot use COPY, it might help to use PREPARE to create a prepared INSERT statement, and then use EXECUTE as many times as required. This avoids some of the overhead of repeatedly parsing and planning INSERT. Different interfaces provide this facility in different ways; look for “prepared statements” in the interface documentation.

Note that loading a large number of rows using COPY is almost always faster than using INSERT, even if PREPARE is used and multiple insertions are batched into a single transaction.

COPY is fastest when used within the same transaction as an earlier CREATE TABLE or TRUNCATE command. In such cases no WAL needs to be written, because in case of an error, the files containing the newly loaded data will be removed anyway. However, this consideration only applies when wal_level is minimal as all commands must write WAL otherwise.

14.4.3. Remove Indexes

If you are loading a freshly created table, the fastest method is to create the table, bulk load the table’s data using COPY, then create any indexes needed for the table. Creating an index on pre-existing data is quicker than updating it incrementally as each row is loaded.

If you are adding large amounts of data to an existing table, it might be a win to drop the indexes, load the table, and then recreate the indexes. Of course, the database performance for other users
might suffer during the time the indexes are missing. One should also think twice before dropping a
unique index, since the error checking afforded by the unique constraint will be lost while the index
is missing.

14.4.4. Remove Foreign Key Constraints

Just as with indexes, a foreign key constraint can be checked “in bulk” more efficiently than row-by-
row. So it might be useful to drop foreign key constraints, load data, and re-create the constraints.
Again, there is a trade-off between data load speed and loss of error checking while the constraint is
missing.

What’s more, when you load data into a table with existing foreign key constraints, each new row
requires an entry in the server’s list of pending trigger events (since it is the firing of a trigger that
checks the row’s foreign key constraint). Loading many millions of rows can cause the trigger event
queue to overflow available memory, leading to intolerable swapping or even outright failure of the
command. Therefore it may be necessary, not just desirable, to drop and re-apply foreign keys when
loading large amounts of data. If temporarily removing the constraint isn’t acceptable, the only other
recourse may be to split up the load operation into smaller transactions.

14.4.5. Increase maintenance_work_mem

Temporarily increasing the maintenance_work_mem configuration variable when loading large
amounts of data can lead to improved performance. This will help to speed up CREATE INDEX
commands and ALTER TABLE ADD FOREIGN KEY commands. It won’t do much for COPY itself, so
this advice is only useful when you are using one or both of the above techniques.

14.4.6. Increase max_wal_size

Temporarily increasing the max_wal_size configuration variable can also make large data loads faster.
This is because loading a large amount of data into PostgreSQL will cause checkpoints to occur
more often than the normal checkpoint frequency (specified by the checkpoint_timeout configu-
ration variable). Whenever a checkpoint occurs, all dirty pages must be flushed to disk. By increasing
max_wal_size temporarily during bulk data loads, the number of checkpoints that are required can
be reduced.

14.4.7. Disable WAL Archival and Streaming Replication

When loading large amounts of data into an installation that uses WAL archiving or streaming repi-
cation, it might be faster to take a new base backup after the load has completed than to process
a large amount of incremental WAL data. To prevent incremental WAL logging while loading, dis-
able archiving and streaming replication, by setting wal_level to minimal, archive_mode to off, and
max_wal_senders to zero. But note that changing these settings requires a server restart.

Aside from avoiding the time for the archiver or WAL sender to process the WAL data, doing this
will actually make certain commands faster, because they are designed not to write WAL at all if
wal_level is minimal. (They can guarantee crash safety more cheaply by doing an fsync at the
end than by writing WAL.) This applies to the following commands:

- CREATE TABLE AS SELECT
Chapter 14. Performance Tips

• CREATE INDEX (and variants such as ALTER TABLE ADD PRIMARY KEY)
• ALTER TABLE SET TABLESPACE
• CLUSTER
• COPY FROM, when the target table has been created or truncated earlier in the same transaction

14.4.8. Run ANALYZE Afterwards

Whenever you have significantly altered the distribution of data within a table, running ANALYZE is strongly recommended. This includes bulk loading large amounts of data into the table. Running ANALYZE (or VACUUM ANALYZE) ensures that the planner has up-to-date statistics about the table. With no statistics or obsolete statistics, the planner might make poor decisions during query planning, leading to poor performance on any tables with inaccurate or nonexistent statistics. Note that if the autovacuum daemon is enabled, it might run ANALYZE automatically; see Section 24.1.3 and Section 24.1.6 for more information.

14.4.9. Some Notes About pg_dump

Dump scripts generated by pg_dump automatically apply several, but not all, of the above guidelines. To reload a pg_dump dump as quickly as possible, you need to do a few extra things manually. (Note that these points apply while restoring a dump, not while creating it. The same points apply whether loading a text dump with psql or using pg_restore to load from a pg_dump archive file.)

By default, pg_dump uses COPY, and when it is generating a complete schema-and-data dump, it is careful to load data before creating indexes and foreign keys. So in this case several guidelines are handled automatically. What is left for you to do is to:

• Set appropriate (i.e., larger than normal) values for maintenance_work_mem and max_wal_size.
• If using WAL archiving or streaming replication, consider disabling them during the restore. To do that, set archive_mode to off, wal_level to minimal, and max_wal_senders to zero before loading the dump. Afterwards, set them back to the right values and take a fresh base backup.
• Experiment with the parallel dump and restore modes of both pg_dump and pg_restore and find the optimal number of concurrent jobs to use. Dumping and restoring in parallel by means of the -j option should give you a significantly higher performance over the serial mode.
• Consider whether the whole dump should be restored as a single transaction. To do that, pass the -1 or --single-transaction command-line option to psql or pg_restore. When using this mode, even the smallest of errors will rollback the entire restore, possibly discarding many hours of processing. Depending on how interrelated the data is, that might seem preferable to manual cleanup, or not. COPY commands will run fastest if you use a single transaction and have WAL archiving turned off.
• If multiple CPUs are available in the database server, consider using pg_restore’s --jobs option. This allows concurrent data loading and index creation.
• Run ANALYZE afterwards.
Chapter 14. Performance Tips

A data-only dump will still use COPY, but it does not drop or recreate indexes, and it does not normally touch foreign keys. ¹ So when loading a data-only dump, it is up to you to drop and recreate indexes and foreign keys if you wish to use those techniques. It’s still useful to increase max_wal_size while loading the data, but don’t bother increasing maintenance_work_mem; rather, you’d do that while manually recreating indexes and foreign keys afterwards. And don’t forget to ANALYZE when you’re done; see Section 24.1.3 and Section 24.1.6 for more information.

14.5. Non-Durable Settings

Durability is a database feature that guarantees the recording of committed transactions even if the server crashes or loses power. However, durability adds significant database overhead, so if your site does not require such a guarantee, PostgreSQL can be configured to run much faster. The following are configuration changes you can make to improve performance in such cases. Except as noted below, durability is still guaranteed in case of a crash of the database software; only abrupt operating system stoppage creates a risk of data loss or corruption when these settings are used.

- Place the database cluster’s data directory in a memory-backed file system (i.e. RAM disk). This eliminates all database disk I/O, but limits data storage to the amount of available memory (and perhaps swap).
- Turn off fsync; there is no need to flush data to disk.
- Turn off synchronous_commit; there might be no need to force WAL writes to disk on every commit. This setting does risk transaction loss (though not data corruption) in case of a crash of the database.
- Turn off full_page_writes; there is no need to guard against partial page writes.
- Increase max_wal_size and checkpoint_timeout; this reduces the frequency of checkpoints, but increases the storage requirements of /pg_xlog.
- Create unlogged tables to avoid WAL writes, though it makes the tables non-crash-safe.

¹ You can get the effect of disabling foreign keys by using the --disable-triggers option — but realize that that eliminates, rather than just postpones, foreign key validation, and so it is possible to insert bad data if you use it.
Chapter 15. Parallel Query

PostgreSQL can devise query plans which can leverage multiple CPUs in order to answer queries faster. This feature is known as parallel query. Many queries cannot benefit from parallel query, either due to limitations of the current implementation or because there is no imaginable query plan which is any faster than the serial query plan. However, for queries that can benefit, the speedup from parallel query is often very significant. Many queries can run more than twice as fast when using parallel query, and some queries can run four times faster or even more. Queries that touch a large amount of data but return only a few rows to the user will typically benefit most. This chapter explains some details of how parallel query works and in which situations it can be used so that users who wish to make use of it can understand what to expect.

15.1. How Parallel Query Works

When the optimizer determines that parallel query is the fastest execution strategy for a particular query, it will create a query plan which includes a **Gather node**. Here is a simple example:

```
EXPLAIN SELECT * FROM pgbench_accounts WHERE filler LIKE '%x%';
```

```
QUERY PLAN
-------------------------------------------------------------------------------------
| Gather (cost=1000.00..217018.43 rows=1 width=97) |
| Workers Planned: 2                                |
| -> Parallel Seq Scan on pgbench_accounts (cost=0.00..216018.33 rows=1 width=97) |
| Filter: (filler ~~ '%x%'::text)                   |
```

In all cases, the **Gather node** will have exactly one child plan, which is the portion of the plan that will be executed in parallel. If the **Gather node** is at the very top of the plan tree, then the entire query will execute in parallel. If it is somewhere else in the plan tree, then only that portion of the query will run in parallel. In the example above, the query accesses only one table, so there is only one plan node other than the **Gather node** itself; since that plan node is a child of the **Gather node**, it will run in parallel.

Using EXPLAIN, you can see the number of workers chosen by the planner. When the **Gather node** is reached during query execution, the process which is implementing the user’s session will request a number of background worker processes equal to the number of workers chosen by the planner. The total number of background workers that can exist at any one time is limited by `max_worker_processes`, so it is possible for a parallel query to run with fewer workers than planned, or even with no workers at all. The optimal plan may depend on the number of workers that are available, so this can result in poor query performance. If this occurrence is frequent, considering increasing `max_worker_processes` so that more workers can be run simultaneously or alternatively reducing `max_parallel_workers_per_gather` so that the planner requests fewer workers.

Every background worker process which is successfully started for a given parallel query will execute the portion of the plan which is a descendent of the **Gather node**. The leader will also execute that portion of the plan, but it has an additional responsibility: it must also read all of the tuples generated by the workers. When the parallel portion of the plan generates only a small number of tuples, the leader will often behave very much like an additional worker, speeding up query execution. Conversely, when the parallel portion of the plan generates a large number of tuples, the leader may be almost entirely occupied with reading the tuples generated by the workers and performing any further
processing steps which are required by plan nodes above the level of the Gather node. In such cases, the leader will do very little of the work of executing the parallel portion of the plan.

## 15.2. When Can Parallel Query Be Used?

There are several settings which can cause the query planner not to generate a parallel query plan under any circumstances. In order for any parallel query plans whatsoever to be generated, the following settings must be configured as indicated.

- `max_parallel_workers_per_gather` must be set to a value which is greater than zero. This is a special case of the more general principle that no more workers should be used than the number configured via `max_parallel_workers_per_gather`.

- `dynamic_shared_memory_type` must be set to a value other than `none`. Parallel query requires dynamic shared memory in order to pass data between cooperating processes.

In addition, the system must not be running in single-user mode. Since the entire database system is running in single process in this situation, no background workers will be available.

Even when it is in general possible for parallel query plans to be generated, the planner will not generate them for a given query if any of the following are true:

- The query writes any data or locks any database rows. If a query contains a data-modifying operation either at the top level or within a CTE, no parallel plans for that query will be generated. This is a limitation of the current implementation which could be lifted in a future release.

- The query might be suspended during execution. In any situation in which the system thinks that partial or incremental execution might occur, no parallel plan is generated. For example, a cursor created using `DECLARE CURSOR` will never use a parallel plan. Similarly, a PL/pgsql loop of the form `FOR x IN query LOOP .. END LOOP` will never use a parallel plan, because the parallel query system is unable to verify that the code in the loop is safe to execute while parallel query is active.

- The query uses any function marked `PARALLEL UNSAFE`. Most system-defined functions are `PARALLEL SAFE`, but user-defined functions are marked `PARALLEL UNSAFE` by default. See the discussion of Section 15.4.

- The query is running inside of another query that is already parallel. For example, if a function called by a parallel query issues an SQL query itself, that query will never use a parallel plan. This is a limitation of the current implementation, but it may not be desirable to remove this limitation, since it could result in a single query using a very large number of processes.

- The transaction isolation level is serializable. This is a limitation of the current implementation.

Even when parallel query plan is generated for a particular query, there are several circumstances under which it will be impossible to execute that plan in parallel at execution time. If this occurs, the leader will execute the portion of the plan below the Gather node entirely by itself, almost as if the Gather node were not present. This will happen if any of the following conditions are met:

- No background workers can be obtained because of the limitation that the total number of background workers cannot exceed `max_worker_processes`.

- The client sends an Execute message with a non-zero fetch count. See the discussion of the extended query protocol. Since `libpq` currently provides no way to send such a message, this can only occur when using a client that does not rely on `libpq`. If this is a frequent occurrence, it may be a good
idea to set max_parallel_workers_per_gather in sessions where it is likely, so as to avoid generating query plans that may be suboptimal when run serially.

• A prepared statement is executed using a `CREATE TABLE .. AS EXECUTE ..` statement. This construct converts what otherwise would have been a read-only operation into a read-write operation, making it ineligible for parallel query.

• The transaction isolation level is serializable. This situation does not normally arise, because parallel query plans are not generated when the transaction isolation level is serializable. However, it can happen if the transaction isolation level is changed to serializable after the plan is generated and before it is executed.

### 15.3. Parallel Plans

Because each worker executes the parallel portion of the plan to completion, it is not possible to simply take an ordinary query plan and run it using multiple workers. Each worker would produce a full copy of the output result set, so the query would not run any faster than normal but would produce incorrect results. Instead, the parallel portion of the plan must be what is known internally to the query optimizer as a partial plan; that is, it must be constructed so that each process which executes the plan will generate only a subset of the output rows in such a way that each required output row is guaranteed to be generated by exactly one of the cooperating processes.

#### 15.3.1. Parallel Scans

Currently, the only type of scan which has been modified to work with parallel query is a sequential scan. Therefore, the driving table in a parallel plan will always be scanned using a `Parallel Seq Scan`. The relation’s blocks will be divided among the cooperating processes. Blocks are handed out one at a time, so that access to the relation remains sequential. Each process will visit every tuple on the page assigned to it before requesting a new page.

#### 15.3.2. Parallel Joins

The driving table may be joined to one or more other tables using nested loops or hash joins. The inner side of the join may be any kind of non-parallel plan that is otherwise supported by the planner provided that it is safe to run within a parallel worker. For example, it may be an index scan which looks up a value taken from the outer side of the join. Each worker will execute the inner side of the join in full, which for hash join means that an identical hash table is built in each worker process.

#### 15.3.3. Parallel Aggregation

PostgreSQL supports parallel aggregation by aggregating in two stages. First, each process participating in the parallel portion of the query performs an aggregation step, producing a partial result for each group of which that process is aware. This is reflected in the plan as a `Partial Aggregate` node. Second, the partial results are transferred to the leader via the `Gather` node. Finally, the leader re-aggregates the results across all workers in order to produce the final result. This is reflected in the plan as a `Finalize Aggregate` node.

Because the `Finalize Aggregate` node runs on the leader process, queries which produce a relatively large number of groups in comparison to the number of input rows will appear less favor-
able to the query planner. For example, in the worst-case scenario the number of groups seen by the Finalize Aggregate node could be as many as the number of input rows which were seen by all worker processes in the Partial Aggregate stage. For such cases, there is clearly going to be no performance benefit to using parallel aggregation. The query planner takes this into account during the planning process and is unlikely to choose parallel aggregate in this scenario.

Parallel aggregation is not supported in all situations. Each aggregate must be safe for parallelism and must have a combine function. If the aggregate has a transition state of type internal, it must have serialization and deserialization functions. See CREATE AGGREGATE for more details. Parallel aggregation is not supported if any aggregate function call contains DISTINCT or ORDER BY clause and is also not supported for ordered set aggregates or when the query involves GROUPING SETS. It can only be used when all joins involved in the query are also part of the parallel portion of the plan.

15.3.4. Parallel Plan Tips

If a query that is expected to do so does not produce a parallel plan, you can try reducing parallel_setup_cost or parallel_tuple_cost. Of course, this plan may turn out to be slower than the serial plan which the planner preferred, but this will not always be the case. If you don’t get a parallel plan even with very small values of these settings (e.g. after setting them both to zero), there may be some reason why the query planner is unable to generate a parallel plan for your query. See Section 15.2 and Section 15.4 for information on why this may be the case.

When executing a parallel plan, you can use EXPLAIN (ANALYZE, VERBOSE) to display per-worker statistics for each plan node. This may be useful in determining whether the work is being evenly distributed between all plan nodes and more generally in understanding the performance characteristics of the plan.

15.4. Parallel Safety

The planner classifies operations involved in a query as either parallel safe, parallel restricted, or parallel unsafe. A parallel safe operation is one which does not conflict with the use of parallel query. A parallel restricted operation is one which cannot be performed in a parallel worker, but which can be performed in the leader while parallel query is in use. Therefore, parallel restricted operations can never occur below a Gather node, but can occur elsewhere in a plan which contains a Gather node. A parallel unsafe operation is one which cannot be performed while parallel query is in use, not even in the leader. When a query contains anything which is parallel unsafe, parallel query is completely disabled for that query.

The following operations are always parallel restricted.

- Scans of common table expressions (CTEs).
- Scans of temporary tables.
- Scans of foreign tables, unless the foreign data wrapper has an IsForeignScanParallelSafe API which indicates otherwise.
- Access to an InitPlan or SubPlan.
15.4.1. Parallel Labeling for Functions and Aggregates

The planner cannot automatically determine whether a user-defined function or aggregate is parallel safe, parallel restricted, or parallel unsafe, because this would require predicting every operation which the function could possibly perform. In general, this is equivalent to the Halting Problem and therefore impossible. Even for simple functions where it could conceivably be done, we do not try, since this would be expensive and error-prone. Instead, all user-defined functions are assumed to be parallel unsafe unless otherwise marked. When using CREATE FUNCTION or ALTER FUNCTION, markings can be set by specifying PARALLEL SAFE, PARALLEL RESTRICTED, or PARALLEL UNSAFE as appropriate. When using CREATE AGGREGATE, the PARALLEL option can be specified with SAFE, RESTRICTED, or UNSAFE as the corresponding value.

Functions and aggregates must be marked PARALLEL UNSAFE if they write to the database, access sequences, change the transaction state even temporarily (e.g. a PL/pgsql function which establishes an EXCEPTION block to catch errors), or make persistent changes to settings. Similarly, functions must be marked PARALLEL RESTRICTED if they access temporary tables, client connection state, cursors, prepared statements, or miscellaneous backend-local state which the system cannot synchronize across workers. For example, setseed and random are parallel restricted for this last reason.

In general, if a function is labeled as being safe when it is restricted or unsafe, or if it is labeled as being restricted when it is in fact unsafe, it may throw errors or produce wrong answers when used in a parallel query. C-language functions could in theory exhibit totally undefined behavior if mislabeled, since there is no way for the system to protect itself against arbitrary C code, but in most likely cases the result will be no worse than for any other function. If in doubt, it is probably best to label functions as UNSAFE.

If a function executed within a parallel worker acquires locks which are not held by the leader, for example by querying a table not referenced in the query, those locks will be released at worker exit, not end of transaction. If you write a function which does this, and this behavior difference is important to you, mark such functions as PARALLEL RESTRICTED to ensure that they execute only in the leader.

Note that the query planner does not consider deferring the evaluation of parallel-restricted functions or aggregates involved in the query in order to obtain a superior plan. So, for example, if a WHERE clause applied to a particular table is parallel restricted, the query planner will not consider placing the scan of that table below a Gather node. In some cases, it would be possible (and perhaps even efficient) to include the scan of that table in the parallel portion of the query and defer the evaluation of the WHERE clause so that it happens above the Gather node. However, the planner does not do this.
III. Server Administration

This part covers topics that are of interest to a PostgreSQL database administrator. This includes installation of the software, set up and configuration of the server, management of users and databases, and maintenance tasks. Anyone who runs a PostgreSQL server, even for personal use, but especially in production, should be familiar with the topics covered in this part.

The information in this part is arranged approximately in the order in which a new user should read it. But the chapters are self-contained and can be read individually as desired. The information in this part is presented in a narrative fashion in topical units. Readers looking for a complete description of a particular command should see Part VI.

The first few chapters are written so they can be understood without prerequisite knowledge, so new users who need to set up their own server can begin their exploration with this part. The rest of this part is about tuning and management; that material assumes that the reader is familiar with the general use of the PostgreSQL database system. Readers are encouraged to look at Part I and Part II for additional information.
Chapter 16. Installation from Source Code

This chapter describes the installation of PostgreSQL using the source code distribution. (If you are installing a pre-packaged distribution, such as an RPM or Debian package, ignore this chapter and read the packager’s instructions instead.)

16.1. Short Version

./configure
make
su
make install
adduser postgres
mkdir /usr/local/pgsql/data
chown postgres /usr/local/pgsql/data
su - postgres
/usr/local/pgsql/bin/initdb -D /usr/local/pgsql/data
/usr/local/pgsql/bin/postgres -D /usr/local/pgsql/data >logfile 2>&1 &
/usr/local/pgsql/bin/createdb test
/usr/local/pgsql/bin/psql test

The long version is the rest of this chapter.

16.2. Requirements

In general, a modern Unix-compatible platform should be able to run PostgreSQL. The platforms that had received specific testing at the time of release are listed in Section 16.6 below. In the doc subdirectory of the distribution there are several platform-specific FAQ documents you might wish to consult if you are having trouble.

The following software packages are required for building PostgreSQL:

- GNU make version 3.80 or newer is required; other make programs or older GNU make versions will not work. (GNU make is sometimes installed under the name gmake.) To test for GNU make enter:

  make --version

- You need an ISO/ANSI C compiler (at least C89-compliant). Recent versions of GCC are re-commended, but PostgreSQL is known to build using a wide variety of compilers from different vendors.

- tar is required to unpack the source distribution, in addition to either gzip or bzip2.

- The GNU Readline library is used by default. It allows psql (the PostgreSQL command line SQL interpreter) to remember each command you type, and allows you to use arrow keys to re-call and edit previous commands. This is very helpful and is strongly recommended. If you don’t want to use it then you must specify the --without-readline option to configure. As an al-ternative, you can often use the BSD-licensed libedit library, originally developed on NetBSD. The libedit library is GNU Readline-compatible and is used if libreadline is not found, or
if --with-libedit-preferred is used as an option to configure. If you are using a package-based Linux distribution, be aware that you need both the readline and readline-devel packages, if those are separate in your distribution.

- The zlib compression library is used by default. If you don’t want to use it then you must specify the --without-zlib option to configure. Using this option disables support for compressed archives in pg_dump and pg_restore.

The following packages are optional. They are not required in the default configuration, but they are needed when certain build options are enabled, as explained below:

- To build the server programming language PL/Perl you need a full Perl installation, including the libperl library and the header files. The minimum required version is Perl 5.8.3. Since PL/Perl will be a shared library, the libperl library must be a shared library also on most platforms. This appears to be the default in recent Perl versions, but it was not in earlier versions, and in any case it is the choice of whomever installed Perl at your site. configure will fail if building PL/Perl is selected but it cannot find a shared libperl. In that case, you will have to rebuild and install Perl manually to be able to build PL/Perl. During the configuration process for Perl, request a shared library.

If you intend to make more than incidental use of PL/Perl, you should ensure that the Perl installation was built with the usemultiplicity option enabled (perl -V will show whether this is the case).

- To build the PL/Python server programming language, you need a Python installation with the header files and the distutils module. The minimum required version is Python 2.3. (To work with function arguments of type numeric, a 2.3.x installation must include the separately-available cdecimal module; note the PL/Python regression tests will not pass if that is missing.) Python 3 is supported if it’s version 3.1 or later; but see Section 44.1 when using Python 3.

Since PL/Python will be a shared library, the libpython library must be a shared library also on most platforms. This is not the case in a default Python installation built from source, but a shared library is available in many operating system distributions. configure will fail if building PL/Python is selected but it cannot find a shared libpython. That might mean that you either have to install additional packages or rebuild (part of) your Python installation to provide this shared library. When building from source, run Python’s configure with the --enable-shared flag.

- To build the PL/Tcl procedural language, you of course need a Tcl installation. The minimum required version is Tcl 8.4.

- To enable Native Language Support (NLS), that is, the ability to display a program’s messages in a language other than English, you need an implementation of the Gettext API. Some operating systems have this built-in (e.g., Linux, NetBSD, Solaris), for other systems you can download an add-on package from http://www.gnu.org/software/gettext/. If you are using the Gettext implementation in the GNU C library then you will additionally need the GNU Gettext package for some utility programs. For any of the other implementations you will not need it.

- You need Kerberos, OpenSSL, OpenLDAP, and/or PAM, if you want to support authentication or encryption using those services.

- To build the PostgreSQL documentation, there is a separate set of requirements; see Section J.2.

If you are building from a Git tree instead of using a released source package, or if you want to do server development, you also need the following packages:
Chapter 16. Installation from Source Code

- GNU Flex and Bison are needed to build from a Git checkout, or if you changed the actual scanner and parser definition files. If you need them, be sure to get Flex 2.5.31 or later and Bison 1.875 or later. Other lex and yacc programs cannot be used.
- Perl 5.8.3 or later is needed to build from a Git checkout, or if you changed the input files for any of the build steps that use Perl scripts. If building on Windows you will need Perl in any case. Perl is also required to run some test suites.

If you need to get a GNU package, you can find it at your local GNU mirror site (see http://www.gnu.org/order/ftp.html for a list) or at ftp://ftp.gnu.org/gnu/.

Also check that you have sufficient disk space. You will need about 100 MB for the source tree during compilation and about 20 MB for the installation directory. An empty database cluster takes about 35 MB; databases take about five times the amount of space that a flat text file with the same data would take. If you are going to run the regression tests you will temporarily need up to an extra 150 MB. Use the df command to check free disk space.

16.3. Getting The Source

The PostgreSQL 9.6.17 sources can be obtained from the download section of our website: https://www.postgresql.org/download/. You should get a file named postgresql-9.6.17.tar.gz or postgresql-9.6.17.tar.bz2. After you have obtained the file, unpack it:

```
gunzip postgresql-9.6.17.tar.gz
tar xf postgresql-9.6.17.tar
```
(Use bunzip2 instead of gunzip if you have the .bz2 file.) This will create a directory postgresql-9.6.17 under the current directory with the PostgreSQL sources. Change into that directory for the rest of the installation procedure.

You can also get the source directly from the version control repository, see Appendix I.

16.4. Installation Procedure

1. Configuration

   The first step of the installation procedure is to configure the source tree for your system and choose the options you would like. This is done by running the configure script. For a default installation simply enter:

   ```
   ./configure
   ```

   This script will run a number of tests to determine values for various system dependent variables and detect any quirks of your operating system, and finally will create several files in the build tree to record what it found. You can also run configure in a directory outside the source tree, if you want to keep the build directory separate. This procedure is also called a VPATH build. Here’s how:

   ```
   mkdir build_dir
   cd build_dir
   /path/to/source/tree/configure [options go here]
   make
   ```
The default configuration will build the server and utilities, as well as all client applications and interfaces that require only a C compiler. All files will be installed under /usr/local/pgsql by default.

You can customize the build and installation process by supplying one or more of the following command line options to configure:

`--prefix=PREFIX`

Install all files under the directory `PREFIX` instead of /usr/local/pgsql. The actual files will be installed into various subdirectories; no files will ever be installed directly into the `PREFIX` directory.

If you have special needs, you can also customize the individual subdirectories with the following options. However, if you leave these with their defaults, the installation will be relocatable, meaning you can move the directory after installation. (The `man` and `doc` locations are not affected by this.)

For relocatable installs, you might want to use `configure`'s `--disable-rpath` option. Also, you will need to tell the operating system how to find the shared libraries.

`--exec-prefix=EXEC-PREFIX`

You can install architecture-dependent files under a different prefix, `EXEC-PREFIX`, than what `PREFIX` was set to. This can be useful to share architecture-independent files between hosts. If you omit this, then `EXEC-PREFIX` is set equal to `PREFIX` and both architecture-dependent and independent files will be installed under the same tree, which is probably what you want.

`--bindir=DIRECTORY`

Specifies the directory for executable programs. The default is `EXEC-PREFIX/bin`, which normally means /usr/local/pgsql/bin.

`--sysconfdir=DIRECTORY`

Sets the directory for various configuration files, `PREFIX/etc` by default.

`--libdir=DIRECTORY`

Sets the location to install libraries and dynamically loadable modules. The default is `EXEC-PREFIX/lib`.

`--includedir=DIRECTORY`

Sets the directory for installing C and C++ header files. The default is `PREFIX/include`.

`--datarootdir=DIRECTORY`

Sets the root directory for various types of read-only data files. This only sets the default for some of the following options. The default is `PREFIX/share`.

`--datadir=DIRECTORY`

Sets the directory for read-only data files used by the installed programs. The default is `DATAROOTDIR`. Note that this has nothing to do with where your database files will be placed.

`--localedir=DIRECTORY`

Sets the directory for installing locale data, in particular message translation catalog files. The default is `DATAROOTDIR/locale`.

459
Chapter 16. Installation from Source Code

--mandir=DIRECTORY

The man pages that come with PostgreSQL will be installed under this directory, in their respective manx subdirectories. The default is DATAROOTDIR/man.

--docdir=DIRECTORY

Sets the root directory for installing documentation files, except “man” pages. This only sets the default for the following options. The default value for this option is DATAROOTDIR/doc/postgresql.

--htmldir=DIRECTORY

The HTML-formatted documentation for PostgreSQL will be installed under this directory. The default is DATAROOTDIR.

Note: Care has been taken to make it possible to install PostgreSQL into shared installation locations (such as /usr/local/include) without interfering with the namespace of the rest of the system. First, the string "postgresql" is automatically appended to datadir, sysconfdir, and docdir, unless the fully expanded directory name already contains the string "postgres" or "pgsql". For example, if you choose /usr/local as prefix, the documentation will be installed in /usr/local/doc/postgresql, but if the prefix is /opt/postgres, then it will be in /opt/postgres/doc. The public C header files of the client interfaces are installed into includedir and are namespace-clean. The internal header files and the server header files are installed into private directories under includedir. See the documentation of each interface for information about how to access its header files. Finally, a private subdirectory will also be created, if appropriate, under libdir for dynamically loadable modules.

--with-extra-version=STRING

Append STRING to the PostgreSQL version number. You can use this, for example, to mark binaries built from unreleased Git snapshots or containing custom patches with an extra version string such as a git describe identifier or a distribution package release number.

--with-includes=DIRECTORIES

DIRECTORY is a colon-separated list of directories that will be added to the list the compiler searches for header files. If you have optional packages (such as GNU Readline) installed in a non-standard location, you have to use this option and probably also the corresponding --with-libraries option.

Example: --with-includes=/opt/gnu/include:/usr/sup/include.

--with-libraries=DIRECTORIES

DIRECTORY is a colon-separated list of directories to search for libraries. You will probably have to use this option (and the corresponding --with-includes option) if you have packages installed in non-standard locations.

Example: --with-libraries=/opt/gnu/lib:/usr/sup/lib.

--enable-nls=[LANGUAGE]

Enables Native Language Support (NLS), that is, the ability to display a program’s messages in a language other than English. LANGUAGES is an optional space-separated list of codes of the languages that you want supported, for example --enable-nls=’de fr’. (The intersection between your list and the set of actually provided translations will be computed automatically.) If you do not specify a list, then all available translations are installed.
Chapter 16. Installation from Source Code

To use this option, you will need an implementation of the Gettext API; see above.

--with-pgport=NUMBER

Set NUMBER as the default port number for server and clients. The default is 5432. The port can always be changed later on, but if you specify it here then both server and clients will have the same default compiled in, which can be very convenient. Usually the only good reason to select a non-default value is if you intend to run multiple PostgreSQL servers on the same machine.

--with-perl

Build the PL/Perl server-side language.

--with-python

Build the PL/Python server-side language.

--with-tcl

Build the PL/Tcl server-side language.

--with-tclconfig=DIRECTORY

Tcl installs the file tclConfig.sh, which contains configuration information needed to build modules interfacing to Tcl. This file is normally found automatically at a well-known location, but if you want to use a different version of Tcl you can specify the directory in which to look for it.

--with-gssapi

Build with support for GSSAPI authentication. On many systems, the GSSAPI (usually a part of the Kerberos installation) system is not installed in a location that is searched by default (e.g., /usr/include, /usr/lib), so you must use the options --with-includes and --with-libraries in addition to this option. configure will check for the required header files and libraries to make sure that your GSSAPI installation is sufficient before proceeding.

--with-krb-srvnam=NAME

The default name of the Kerberos service principal used by GSSAPI postgres is the default. There’s usually no reason to change this unless you have a Windows environment, in which case it must be set to upper case POSTGRES.

--with-openssl

Build with support for SSL (encrypted) connections. This requires the OpenSSL package to be installed. configure will check for the required header files and libraries to make sure that your OpenSSL installation is sufficient before proceeding.

--with-pam

Build with PAM (Pluggable Authentication Modules) support.

--with-bsd-auth

Build with BSD Authentication support. (The BSD Authentication framework is currently only available on OpenBSD.)

--with-ldap

Build with LDAP support for authentication and connection parameter lookup (see Section 32.17 and Section 20.3.7 for more information). On Unix, this requires the OpenLDAP package to be installed. On Windows, the default WinLDAP library is used. configure
Chapter 16. Installation from Source Code

will check for the required header files and libraries to make sure that your OpenLDAP installation is sufficient before proceeding.

--with-systemd

Build with support for systemd service notifications. This improves integration if the server binary is started under systemd but has no impact otherwise; see Section 18.3 for more information. libsystemd and the associated header files need to be installed to be able to use this option.

--without-readline

Prevents use of the Readline library (and libedit as well). This option disables command-line editing and history in psql, so it is not recommended.

--with-libedit-preferred

Favors the use of the BSD-licensed libedit library rather than GPL-licensed Readline. This option is significant only if you have both libraries installed; the default in that case is to use Readline.

--with-bonjour

Build with Bonjour support. This requires Bonjour support in your operating system. Recommended on OS X.

--with-uuid=LIBRARY

Build the uuid-ossp module (which provides functions to generate UUIDs), using the specified UUID library. LIBRARY must be one of:

- bsd to use the UUID functions found in FreeBSD, NetBSD, and some other BSD-derived systems
- e2fs to use the UUID library created by the e2fsprogs project; this library is present in most Linux systems and in OS X, and can be obtained for other platforms as well
- ossp to use the OSSP UUID library\(^1\)

--with-ossp-uuid

Obsolete equivalent of --with-uuid=oss.

--with-libxml

Build with libxml (enables SQL/XML support). Libxml version 2.6.23 or later is required for this feature.

Libxml installs a program xml2-config that can be used to detect the required compiler and linker options. PostgreSQL will use it automatically if found. To specify a libxml installation at an unusual location, you can either set the environment variable XML2_CONFIG to point to the xml2-config program belonging to the installation, or use the options --with-includes and --with-libraries.

--with-libxslt

Use libxslt when building the xml2 module. xml2 relies on this library to perform XSL transformations of XML.

--disable-integer-datetimes

Disable support for 64-bit integer storage for timestamps and intervals, and store datetime values as floating-point numbers instead. Floating-point datetime storage was the default in

\(^1\) http://www.ossp.org/pkg/lib/uuid/
PostgreSQL releases prior to 8.4, but it is now deprecated, because it does not support microsecond precision for the full range of timestamp values. However, integer-based date-time storage requires a 64-bit integer type. Therefore, this option can be used when no such type is available, or for compatibility with applications written for prior versions of PostgreSQL. See Section 8.5 for more information.

--disable-float4-byval
Disable passing float4 values “by value”, causing them to be passed “by reference” instead. This option costs performance, but may be needed for compatibility with old user-defined functions that are written in C and use the “version 0” calling convention. A better long-term solution is to update any such functions to use the “version 1” calling convention.

--disable-float8-byval
Disable passing float8 values “by value”, causing them to be passed “by reference” instead. This option costs performance, but may be needed for compatibility with old user-defined functions that are written in C and use the “version 0” calling convention. A better long-term solution is to update any such functions to use the “version 1” calling convention. Note that this option affects not only float8, but also int8 and some related types such as timestamp. On 32-bit platforms, --disable-float8-byval is the default and it is not allowed to select --enable-float8-byval.

--with-segsize=SEGSIZE
Set the segment size, in gigabytes. Large tables are divided into multiple operating-system files, each of size equal to the segment size. This avoids problems with file size limits that exist on many platforms. The default segment size, 1 gigabyte, is safe on all supported platforms. If your operating system has “largefile” support (which most do, nowadays), you can use a larger segment size. This can be helpful to reduce the number of file descriptors consumed when working with very large tables. But be careful not to select a value larger than is supported by your platform and the file systems you intend to use. Other tools you might wish to use, such as tar, could also set limits on the usable file size. It is recommended, though not absolutely required, that this value be a power of 2. Note that changing this value requires an initdb.

--with-blocksize=BLOCKSIZE
Set the block size, in kilobytes. This is the unit of storage and I/O within tables. The default, 8 kilobytes, is suitable for most situations; but other values may be useful in special cases. The value must be a power of 2 between 1 and 32 (kilobytes). Note that changing this value requires an initdb.

--with-wal-segsize=SEGSIZE
Set the WAL segment size, in megabytes. This is the size of each individual file in the WAL log. It may be useful to adjust this size to control the granularity of WAL log shipping. The default size is 16 megabytes. The value must be a power of 2 between 1 and 64 (megabytes). Note that changing this value requires an initdb.

--with-wal-blocksize=BLOCKSIZE
Set the WAL block size, in kilobytes. This is the unit of storage and I/O within the WAL log. The default, 8 kilobytes, is suitable for most situations; but other values may be useful in special cases. The value must be a power of 2 between 1 and 64 (kilobytes). Note that changing this value requires an initdb.

--disable-spinlocks
Allow the build to succeed even if PostgreSQL has no CPU spinlock support for the plat-
Chapter 16. Installation from Source Code

form. The lack of spinlock support will result in poor performance; therefore, this option should only be used if the build aborts and informs you that the platform lacks spinlock support. If this option is required to build PostgreSQL on your platform, please report the problem to the PostgreSQL developers.

--disable-thread-safety

Disable the thread-safety of client libraries. This prevents concurrent threads in libpq and ECPG programs from safely controlling their private connection handles.

--with-system-tzdata=DIRECTORY

PostgreSQL includes its own time zone database, which it requires for date and time operations. This time zone database is in fact compatible with the IANA time zone database provided by many operating systems such as FreeBSD, Linux, and Solaris, so it would be redundant to install it again. When this option is used, the system-supplied time zone database in DIRECTORY is used instead of the one included in the PostgreSQL source distribution. DIRECTORY must be specified as an absolute path. /usr/share/zoneinfo is a likely directory on some operating systems. Note that the installation routine will not detect mismatching or erroneous time zone data. If you use this option, you are advised to run the regression tests to verify that the time zone data you have pointed to works correctly with PostgreSQL.

This option is mainly aimed at binary package distributors who know their target operating system well. The main advantage of using this option is that the PostgreSQL package won’t need to be upgraded whenever any of the many local daylight-saving time rules change. Another advantage is that PostgreSQL can be cross-compiled more straightforwardly if the time zone database files do not need to be built during the installation.

--without-zlib

Prevents use of the Zlib library. This disables support for compressed archives in pg_dump and pg_restore. This option is only intended for those rare systems where this library is not available.

--enable-debug

Compiles all programs and libraries with debugging symbols. This means that you can run the programs in a debugger to analyze problems. This enlarges the size of the installed executables considerably, and on non-GCC compilers it usually also disables compiler optimization, causing slowdowns. However, having the symbols available is extremely helpful for dealing with any problems that might arise. Currently, this option is recommended for production installations only if you use GCC. But you should always have it on if you are doing development work or running a beta version.

--enable-coverage

If using GCC, all programs and libraries are compiled with code coverage testing instrumentation. When run, they generate files in the build directory with code coverage metrics. See Section 31.5 for more information. This option is for use only with GCC and when doing development work.

--enable-profiling

If using GCC, all programs and libraries are compiled so they can be profiled. On backend exit, a subdirectory will be created that contains the gmon.out file for use in profiling. This option is for use only with GCC and when doing development work.
Chapter 16. Installation from Source Code

--enable-cassert

Enables assertion checks in the server, which test for many “cannot happen” conditions. This is invaluable for code development purposes, but the tests can slow down the server significantly. Also, having the tests turned on won’t necessarily enhance the stability of your server! The assertion checks are not categorized for severity, and so what might be a relatively harmless bug will still lead to server restarts if it triggers an assertion failure. This option is not recommended for production use, but you should have it on for development work or when running a beta version.

--enable-depend

Enables automatic dependency tracking. With this option, the makefiles are set up so that all affected object files will be rebuilt when any header file is changed. This is useful if you are doing development work, but is just wasted overhead if you intend only to compile once and install. At present, this option only works with GCC.

--enable-dtrace

Compiles PostgreSQL with support for the dynamic tracing tool DTrace. See Section 28.5 for more information.

To point to the dtrace program, the environment variable DTRACE can be set. This will often be necessary because dtrace is typically installed under /usr/sbin, which might not be in the path.

Extra command-line options for the dtrace program can be specified in the environment variable DTRACEFLAGS. On Solaris, to include DTrace support in a 64-bit binary, you must specify DTRACEFLAGS="-64" to configure. For example, using the GCC compiler:

./configure CC='gcc -m64' --enable-dtrace DTRACEFLAGS='-64' ...

Using Sun’s compiler:

./configure CC='/opt/SUNWpro/bin/cc -xtarget=native64' --enable-dtrace DTRACEFLAGS='-64' ...

--enable-tap-tests

Enable tests using the Perl TAP tools. This requires a Perl installation and the Perl module IPC::Run. See Section 31.4 for more information.

If you prefer a C compiler different from the one configure picks, you can set the environment variable CC to the program of your choice. By default, configure will pick gcc if available, else the platform’s default (usually cc). Similarly, you can override the default compiler flags if needed with the CFLAGS variable.

You can specify environment variables on the configure command line, for example:

./configure CC=/opt/bin/gcc CFLAGS='"-O2 -pipe"

Here is a list of the significant variables that can be set in this manner:

BISON

Bison program

CC

C compiler

CFLAGS

options to pass to the C compiler
Chapter 16. Installation from Source Code

CPP
C preprocessor

CPPFLAGS
options to pass to the C preprocessor

DTRACE
location of the dtrace program

DTRACEFLAGS
options to pass to the dtrace program

FLEX
Flex program

LDFLAGS
options to use when linking either executables or shared libraries

LDFLAGS_EX
additional options for linking executables only

LDFLAGS_SL
additional options for linking shared libraries only

MSGFMT
msgfmt program for native language support

PERL
Full path name of the Perl interpreter. This will be used to determine the dependencies for building PL/Perl.

PYTHON
Full path name of the Python interpreter. This will be used to determine the dependencies for building PL/Python. Also, whether Python 2 or 3 is specified here (or otherwise implicitly chosen) determines which variant of the PL/Python language becomes available. See Section 44.1 for more information. If this is not set, the following are probed in this order: python python3 python2.

TCLSH
Full path name of the Tcl interpreter. This will be used to determine the dependencies for building PL/Tcl, and it will be substituted into Tcl scripts.

XML2_CONFIG
xml2-config program used to locate the libxml installation.

Sometimes it is useful to add compiler flags after-the-fact to the set that were chosen by configure. An important example is that gcc’s -Werror option cannot be included in the CFLAGS passed to configure, because it will break many of configure’s built-in tests. To add such flags, include them in the COPT environment variable while running make. The contents of COPT are added to both the CFLAGS and LDFLAGS options set up by configure. For example, you could do

make COPT=''-Werror''
or
export COPT='--Werror'
make

Note: When developing code inside the server, it is recommended to use the configure options --enable-cassert (which turns on many run-time error checks) and --enable-debug (which improves the usefulness of debugging tools).

If using GCC, it is best to build with an optimization level of at least -O1, because using no optimization (-O0) disables some important compiler warnings (such as the use of uninitialized variables). However, non-zero optimization levels can complicate debugging because stepping through compiled code will usually not match up one-to-one with source code lines. If you get confused while trying to debug optimized code, recompile the specific files of interest with -O0. An easy way to do this is by passing an option to make: make PROFILE=-O0 file.o.

The COPT and PROFILE environment variables are actually handled identically by the PostgreSQL makefiles. Which to use is a matter of preference, but a common habit among developers is to use PROFILE for one-time flag adjustments, while COPT might be kept set all the time.

2. Build

To start the build, type:

make

(Remember to use GNU make.) The build will take a few minutes depending on your hardware. The last line displayed should be:

All of PostgreSQL successfully made. Ready to install.

If you want to build everything that can be built, including the documentation (HTML and man pages), and the additional modules (contrib), type instead:

make world

The last line displayed should be:

PostgreSQL, contrib, and documentation successfully made. Ready to install.

3. Regression Tests

If you want to test the newly built server before you install it, you can run the regression tests at this point. The regression tests are a test suite to verify that PostgreSQL runs on your machine in the way the developers expected it to. Type:

make check

(This won’t work as root; do it as an unprivileged user.) Chapter 31 contains detailed information about interpreting the test results. You can repeat this test at any later time by issuing the same command.

4. Installing the Files

Note: If you are upgrading an existing system be sure to read Section 18.6 which has instructions about upgrading a cluster.

To install PostgreSQL enter:

make install

This will install files into the directories that were specified in step 1. Make sure that you have appropriate permissions to write into that area. Normally you need to do this step as root. Alter-
natively, you can create the target directories in advance and arrange for appropriate permissions to be granted.

To install the documentation (HTML and man pages), enter:

```make install-docs```

If you built the world above, type instead:

```make install-world```

This also installs the documentation.

You can use `make install-strip` instead of `make install` to strip the executable files and libraries as they are installed. This will save some space. If you built with debugging support, stripping will effectively remove the debugging support, so it should only be done if debugging is no longer needed. `install-strip` tries to do a reasonable job saving space, but it does not have perfect knowledge of how to strip every unneeded byte from an executable file, so if you want to save all the disk space you possibly can, you will have to do manual work.

The standard installation provides all the header files needed for client application development as well as for server-side program development, such as custom functions or data types written in C. (Prior to PostgreSQL 8.0, a separate `make install-all-headers` command was needed for the latter, but this step has been folded into the standard install.)

**Client-only installation:** If you want to install only the client applications and interface libraries, then you can use these commands:

```make -C src/bin install```
```make -C src/include install```
```make -C src/interfaces install```
```make -C doc install```

`src/bin` has a few binaries for server-only use, but they are small.

**Uninstallation:** To undo the installation use the command `make uninstall`. However, this will not remove any created directories.

**Cleaning:** After the installation you can free disk space by removing the built files from the source tree with the command `make clean`. This will preserve the files made by the configure program, so that you can rebuild everything with `make` later on. To reset the source tree to the state in which it was distributed, use `make distclean`. If you are going to build for several platforms within the same source tree you must do this and re-configure for each platform. (Alternatively, use a separate build tree for each platform, so that the source tree remains unmodified.)

If you perform a build and then discover that your `configure` options were wrong, or if you change anything that `configure` investigates (for example, software upgrades), then it’s a good idea to do `make distclean` before reconfiguring and rebuilding. Without this, your changes in configuration choices might not propagate everywhere they need to.

### 16.5. Post-Installation Setup

#### 16.5.1. Shared Libraries

On some systems with shared libraries you need to tell the system how to find the newly installed shared libraries. The systems on which this is *not* necessary include FreeBSD, HP-UX, Linux, NetBSD, OpenBSD, and Solaris.
Chapter 16. Installation from Source Code

The method to set the shared library search path varies between platforms, but the most widely-used method is to set the environment variable `LD_LIBRARY_PATH` like so: In Bourne shells (`sh`, `ksh`, `bash`, `zsh`):

```bash
LD_LIBRARY_PATH=/usr/local/pgsql/lib
export LD_LIBRARY_PATH
```

or in `csh` or `tcsh`:

```csh
setenv LD_LIBRARY_PATH /usr/local/pgsql/lib
```

Replace `/usr/local/pgsql/lib` with whatever you set `--libdir` to in step 1. You should put these commands into a shell start-up file such as `/etc/profile` or `~/.bash_profile`. Some good information about the caveats associated with this method can be found at http://xahlee.org/UnixResource_dir/_/ldpath.html.

On some systems it might be preferable to set the environment variable `LD_RUN_PATH` before building.

On Cygwin, put the library directory in the `PATH` or move the `.dll` files into the `bin` directory.

If in doubt, refer to the manual pages of your system (perhaps `ld.so` or `rld`). If you later get a message like:

```bash
psql: error in loading shared libraries
libpq.so.2.1: cannot open shared object file: No such file or directory
```

then this step was necessary. Simply take care of it then.

If you are on Linux and you have root access, you can run:

```bash
/sbin/ldconfig /usr/local/pgsql/lib
```

(or equivalent directory) after installation to enable the run-time linker to find the shared libraries faster. Refer to the manual page of `ldconfig` for more information. On FreeBSD, NetBSD, and OpenBSD the command is:

```bash
/sbin/ldconfig -m /usr/local/pgsql/lib
```

instead. Other systems are not known to have an equivalent command.

### 16.5.2. Environment Variables

If you installed into `/usr/local/pgsql` or some other location that is not searched for programs by default, you should add `/usr/local/pgsql/bin` (or whatever you set `--bindir` to in step 1) into your `PATH`. Strictly speaking, this is not necessary, but it will make the use of PostgreSQL much more convenient.

To do this, add the following to your shell start-up file, such as `~/bash_profile` (or `~/etc/profile`, if you want it to affect all users):

```bash
PATH=/usr/local/pgsql/bin:$PATH
export PATH
```

If you are using `csh` or `tcsh`, then use this command:

```csh
set path = ( /usr/local/pgsql/bin $path )
```
To enable your system to find the man documentation, you need to add lines like the following to a shell start-up file unless you installed into a location that is searched by default:

```sh
MANPATH=/usr/local/pgsql/share/man:$MANPATH
export MANPATH
```

The environment variables `PGHOST` and `PGPORT` specify to client applications the host and port of the database server, overriding the compiled-in defaults. If you are going to run client applications remotely then it is convenient if every user that plans to use the database sets `PGHOST`. This is not required, however; the settings can be communicated via command line options to most client programs.

### 16.6. Supported Platforms

A platform (that is, a CPU architecture and operating system combination) is considered supported by the PostgreSQL development community if the code contains provisions to work on that platform and it has recently been verified to build and pass its regression tests on that platform. Currently, most testing of platform compatibility is done automatically by test machines in the PostgreSQL Build Farm². If you are interested in using PostgreSQL on a platform that is not represented in the build farm, but on which the code works or can be made to work, you are strongly encouraged to set up a build farm member machine so that continued compatibility can be assured.

In general, PostgreSQL can be expected to work on these CPU architectures: x86, x86_64, IA64, PowerPC, PowerPC 64, S/390, S/390x, Sparc, Sparc 64, ARM, MIPS, MIPSEL, M68K, and PA-RISC. Code support exists for M32R and VAX, but these architectures are not known to have been tested recently. It is often possible to build on an unsupported CPU type by configuring with `--disable-spinlocks`, but performance will be poor.

PostgreSQL can be expected to work on these operating systems: Linux (all recent distributions), Windows (Win2000 SP4 and later), FreeBSD, OpenBSD, NetBSD, OS X, AIX, HP/UX, Solaris, and UnixWare. Other Unix-like systems may also work but are not currently being tested. In most cases, all CPU architectures supported by a given operating system will work. Look in the Section 16.7 below to see if there is information specific to your operating system, particularly if using an older system.

If you have installation problems on a platform that is known to be supported according to recent build farm results, please report it to `<pgsql-bugs@lists.postgresql.org>`. If you are interested in porting PostgreSQL to a new platform, `<pgsql-hackers@lists.postgresql.org>` is the appropriate place to discuss that.

### 16.7. Platform-specific Notes

This section documents additional platform-specific issues regarding the installation and setup of PostgreSQL. Be sure to read the installation instructions, and in particular Section 16.2 as well. Also, check Chapter 31 regarding the interpretation of regression test results.

Platforms that are not covered here have no known platform-specific installation issues.

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² [https://buildfarm.postgresql.org/](https://buildfarm.postgresql.org/)
16.7.1. AIX

PostgreSQL works on AIX, but getting it installed properly can be challenging. AIX versions from 4.3.3 to 6.1 are considered supported. You can use GCC or the native IBM compiler xlc. In general, using recent versions of AIX and PostgreSQL helps. Check the build farm for up to date information about which versions of AIX are known to work.

The minimum recommended fix levels for supported AIX versions are:

AIX 4.3.3
   Maintenance Level 11 + post ML11 bundle
AIX 5.1
   Maintenance Level 9 + post ML9 bundle
AIX 5.2
   Technology Level 10 Service Pack 3
AIX 5.3
   Technology Level 7
AIX 6.1
   Base Level

To check your current fix level, use `oslevel -r` in AIX 4.3.3 to AIX 5.2 ML 7, or `oslevel -s` in later versions.

Use the following `configure` flags in addition to your own if you have installed Readline or libz in /usr/local:

```
--with-includes=/usr/local/include
--with-libraries=/usr/local/lib
```

16.7.1.1. GCC Issues

On AIX 5.3, there have been some problems getting PostgreSQL to compile and run using GCC.

You will want to use a version of GCC subsequent to 3.3.2, particularly if you use a prepackaged version. We had good success with 4.0.1. Problems with earlier versions seem to have more to do with the way IBM packaged GCC than with actual issues with GCC, so that if you compile GCC yourself, you might well have success with an earlier version of GCC.

16.7.1.2. Unix-Domain Sockets Broken

AIX 5.3 has a problem where `sockaddr_storage` is not defined to be large enough. In version 5.3, IBM increased the size of `sockaddr_un`, the address structure for Unix-domain sockets, but did not correspondingly increase the size of `sockaddr_storage`. The result of this is that attempts to use Unix-domain sockets with PostgreSQL lead to `libpq` overflowing the data structure. TCP/IP connections work OK, but not Unix-domain sockets, which prevents the regression tests from working.

The problem was reported to IBM, and is recorded as bug report PMR29657. If you upgrade to maintenance level 5300-03 or later, that will include this fix. A quick workaround is to alter `_SS_MAXSIZE` to 1025 in `/usr/include/sys/socket.h`. In either case, recompile PostgreSQL once you have the corrected header file.
16.7.1.3. Internet Address Issues

PostgreSQL relies on the system’s `getaddrinfo` function to parse IP addresses in `listen_addresses`, `pg_hba.conf`, etc. Older versions of AIX have assorted bugs in this function. If you have problems related to these settings, updating to the appropriate AIX fix level shown above should take care of it.

One user reports:

When implementing PostgreSQL version 8.1 on AIX 5.3, we periodically ran into problems where the statistics collector would “mysteriously” not come up successfully. This appears to be the result of unexpected behavior in the IPv6 implementation. It looks like PostgreSQL and IPv6 do not play very well together on AIX 5.3.

Any of the following actions “fix” the problem.

- Delete the IPv6 address for localhost:
  
  (as root)
  
  `# ifconfig lo0 inet6 ::1/0 delete`

- Remove IPv6 from net services. The file `/etc/netsvc.conf` on AIX is roughly equivalent to `/etc/nsswitch.conf` on Solaris/Linux. The default, on AIX, is thus:

  ```
  hosts=local,bind
  ```

  Replace this with:

  ```
  hosts=local4,bind4
  ```

  to deactivate searching for IPv6 addresses.

**Warning**

This is really a workaround for problems relating to immaturity of IPv6 support, which improved visibly during the course of AIX 5.3 releases. It has worked with AIX version 5.3, but does not represent an elegant solution to the problem. It has been reported that this workaround is not only unnecessary, but causes problems on AIX 6.1, where IPv6 support has become more mature.

16.7.1.4. Memory Management

AIX can be somewhat peculiar with regards to the way it does memory management. You can have a server with many multiples of gigabytes of RAM free, but still get out of memory or address space errors when running applications. One example is `createlang` failing with unusual errors. For example, running as the owner of the PostgreSQL installation:

```
-bash-3.00$ createlang plperl template1
createlang: language installation failed: ERROR: could not load library "/opt/dbs/pgsql748/lib/plperl.so": A memory address is not in the address space for the process.
```

Running as a non-owner in the group possessing the PostgreSQL installation:

```
-bash-3.00$ createlang plperl template1
createlang: language installation failed: ERROR: could not load library "/opt/dbs/pgsql748/lib/plperl.so": Bad address
```

Another example is out of memory errors in the PostgreSQL server logs, with every memory allocation near or greater than 256 MB failing.
The overall cause of all these problems is the default bittedness and memory model used by the server process. By default, all binaries built on AIX are 32-bit. This does not depend upon hardware type or kernel in use. These 32-bit processes are limited to 4 GB of memory laid out in 256 MB segments using one of a few models. The default allows for less than 256 MB in the heap as it shares a single segment with the stack.

In the case of the `createlang` example, above, check your umask and the permissions of the binaries in your PostgreSQL installation. The binaries involved in that example were 32-bit and installed as mode 750 instead of 755. Due to the permissions being set in this fashion, only the owner or a member of the possessing group can load the library. Since it isn’t world-readable, the loader places the object into the process’ heap instead of the shared library segments where it would otherwise be placed.

The “ideal” solution for this is to use a 64-bit build of PostgreSQL, but that is not always practical, because systems with 32-bit processors can build, but not run, 64-bit binaries.

If a 32-bit binary is desired, set `LDR_CNTRL` to `MAXDATA=0x00000000`, where `1 <= n <= 8`, before starting the PostgreSQL server, and try different values and `postgresql.conf` settings to find a configuration that works satisfactorily. This use of `LDR_CNTRL` tells AIX that you want the server to have `MAXDATA` bytes set aside for the heap, allocated in 256 MB segments. When you find a workable configuration, `ldedit` can be used to modify the binaries so that they default to using the desired heap size. PostgreSQL can also be rebuilt, passing `configure LDFLAGS="-Wl,-bmaxdata:0x00000000"` to achieve the same effect.

For a 64-bit build, set `OBJECT_MODE` to 64 and pass `CC="gcc -maix64"` and `LDFLAGS="-Wl,-bbigtoc"` to configure. (Options for `xlc` might differ.) If you omit the export of `OBJECT_MODE`, your build may fail with linker errors. When `OBJECT_MODE` is set, it tells AIX’s build utilities such as `ar`, `as`, and `ld` what type of objects to default to handling.

By default, overcommit of paging space can happen. While we have not seen this occur, AIX will kill processes when it runs out of memory and the overcommit is accessed. The closest to this that we have seen is fork failing because the system decided that there was not enough memory for another process. Like many other parts of AIX, the paging space allocation method and out-of-memory kill is configurable on a system- or process-wide basis if this becomes a problem.

References and Resources

“Large Program Support”, AIX Documentation: General Programming Concepts: Writing and Debugging Programs.

“Program Address Space Overview”, AIX Documentation: General Programming Concepts: Writing and Debugging Programs.


Developing and Porting C and C++ Applications on AIX®, IBM Redbook.

16.7.2. Cygwin

PostgreSQL can be built using Cygwin, a Linux-like environment for Windows, but that method is inferior to the native Windows build (see Chapter 17) and running a server under Cygwin is no longer recommended.

When building from source, proceed according to the normal installation procedure (i.e., ./configure; make; etc.), noting the following-Cygwin specific differences:

- Set your path to use the Cygwin bin directory before the Windows utilities. This will help prevent problems with compilation.
- The adduser command is not supported; use the appropriate user management application on Windows NT, 2000, or XP. Otherwise, skip this step.
- The su command is not supported; use ssh to simulate su on Windows NT, 2000, or XP. Otherwise, skip this step.
- OpenSSL is not supported.
- Start cygserver for shared memory support. To do this, enter the command /usr/sbin/cygserver &. This program needs to be running anytime you start the PostgreSQL server or initialize a database cluster (initdb). The default cygserver configuration may need to be changed (e.g., increase SEMMNS) to prevent PostgreSQL from failing due to a lack of system resources.
- Building might fail on some systems where a locale other than C is in use. To fix this, set the locale to C by doing export LANG=C.utf8 before building, and then setting it back to the previous setting, after you have installed PostgreSQL.
- The parallel regression tests (make check) can generate spurious regression test failures due to overflowing the listen() backlog queue which causes connection refused errors or hangs. You can limit the number of connections using the make variable MAX_CONNECTIONS thus:
  make MAX_CONNECTIONS=5 check
(On some systems you can have up to about 10 simultaneous connections).

It is possible to install cygserver and the PostgreSQL server as Windows NT services. For information on how to do this, please refer to the README document included with the PostgreSQL binary package on Cygwin. It is installed in the directory /usr/share/doc/Cygwin.

16.7.3. HP-UX

PostgreSQL 7.3+ should work on Series 700/800 PA-RISC machines running HP-UX 10.X or 11.X, given appropriate system patch levels and build tools. At least one developer routinely tests on HP-UX 10.20, and we have reports of successful installations on HP-UX 11.00 and 11.11.

Aside from the PostgreSQL source distribution, you will need GNU make (HP's make will not do), and either GCC or HP’s full ANSI C compiler. If you intend to build from Git sources rather than a distribution tarball, you will also need Flex (GNU lex) and Bison (GNU yacc). We also recommend making sure you are fairly up-to-date on HP patches. At a minimum, if you are building 64 bit binaries on HP-UX 11.11 you may need PHSS_30966 (11.11) or a successor patch otherwise initdb may hang:

PHSS_30966 s700_800 ld(1) and linker tools cumulative patch
On general principles you should be current on libc and ld/dld patches, as well as compiler patches if
you are using HP’s C compiler. See HP’s support sites such as http://itrc.hp.com and ftp://us-
ffs.external.hp.com/ for free copies of their latest patches.

If you are building on a PA-RISC 2.0 machine and want to have 64-bit binaries using GCC, you
must use GCC 64-bit version. GCC binaries for HP-UX PA-RISC and Itanium are available from
http://www.hp.com/go/gcc. Don’t forget to get and install binutils at the same time.

If you are building on a PA-RISC 2.0 machine and want the compiled binaries to run on PA-RISC 1.1
machines you will need to specify +DAportable in CFLAGS.

If you are building on a HP-UX Itanium machine, you will need the latest HP ANSI C compiler with
its dependent patch or successor patches:

PHSS_30848 s700_800 HP C Compiler (A.05.57)
PHSS_30849 s700_800 u2comp/be/plugin library Patch

If you have both HP’s C compiler and GCC’s, then you might want to explicitly select the compiler
to use when you run configure:

./configure CC=cc

for HP’s C compiler, or

./configure CC=gcc

for GCC. If you omit this setting, then configure will pick gcc if it has a choice.

The default install target location is /usr/local/pgsql, which you might want to change to some-
thing under /opt. If so, use the --prefix switch to configure.

In the regression tests, there might be some low-order-digit differences in the geometry tests, which
vary depending on which compiler and math library versions you use. Any other error is cause for
suspicion.

16.7.4. macOS

On recent macOS releases, it’s necessary to embed the “sysroot” path in the include switches used to
find some system header files. This results in the outputs of the configure script varying depending on
which SDK version was used during configure. That shouldn’t pose any problem in simple scenarios,
but if you are trying to do something like building an extension on a different machine than the server
code was built on, you may need to force use of a different sysroot path. To do that, set PG_SYSROOT,
for example

make PG_SYSROOT=/desired/path all

To find out the appropriate path on your machine, run

xcodebuild -version -sdk macosx Path

Note that building an extension using a different sysroot version than was used to build the core server
is not really recommended; in the worst case it could result in hard-to-debug ABI inconsistencies.

You can also select a non-default sysroot path when configuring, by specifying PG_SYSROOT to con-
figure:
macOS’s “System Integrity Protection” (SIP) feature breaks make check, because it prevents passing the needed setting of `DYLD_LIBRARY_PATH` down to the executables being tested. You can work around that by doing `make install` before `make check`. Most Postgres developers just turn off SIP, though.

### 16.7.5. MinGW/Native Windows

PostgreSQL for Windows can be built using MinGW, a Unix-like build environment for Microsoft operating systems, or using Microsoft’s Visual C++ compiler suite. The MinGW build variant uses the normal build system described in this chapter; the Visual C++ build works completely differently and is described in Chapter 17. It is a fully native build and uses no additional software like MinGW.

A ready-made installer is available on the main PostgreSQL web site.

The native Windows port requires a 32 or 64-bit version of Windows 2000 or later. Earlier operating systems do not have sufficient infrastructure (but Cygwin may be used on those). MinGW, the Unix-like build tools, and MSYS, a collection of Unix tools required to run shell scripts like configure, can be downloaded from http://www.mingw.org/. Neither is required to run the resulting binaries; they are needed only for creating the binaries.

To build 64 bit binaries using MinGW, install the 64 bit tool set from http://mingw-w64.sourceforge.net/, put its bin directory in the `PATH`, and run `configure` with the `--host=x86_64-w64-mingw32` option.

After you have everything installed, it is suggested that you run psql under `CMD.EXE`, as the MSYS console has buffering issues.

#### 16.7.5.1. Collecting Crash Dumps on Windows

If PostgreSQL on Windows crashes, it has the ability to generate minidumps that can be used to track down the cause for the crash, similar to core dumps on Unix. These dumps can be read using the Windows Debugger Tools or using Visual Studio. To enable the generation of dumps on Windows, create a subdirectory named `crashdumps` inside the cluster data directory. The dumps will then be written into this directory with a unique name based on the identifier of the crashing process and the current time of the crash.

### 16.7.6. SCO OpenServer and SCO UnixWare

PostgreSQL can be built on SCO UnixWare 7 and SCO OpenServer 5. On OpenServer, you can use either the OpenServer Development Kit or the Universal Development Kit. However, some tweaking may be needed, as described below.

#### 16.7.6.1. Skunkware

You should locate your copy of the SCO Skunkware CD. The Skunkware CD is included with UnixWare 7 and current versions of OpenServer 5. Skunkware includes ready-to-install versions of many popular programs that are available on the Internet. For example, gzip, gunzip, GNU Make, Flex, and Bison are all included. For UnixWare 7.1, this CD is now labeled "Open
Chapter 16. Installation from Source Code

License Software Supplement”. If you do not have this CD, the software on it is available from http://www.sco.com/skunkware/.

Skunkware has different versions for UnixWare and OpenServer. Make sure you install the correct version for your operating system, except as noted below.

On UnixWare 7.1.3 and beyond, the GCC compiler is included on the UDK CD as is GNU Make.

16.7.6.2. GNU Make

You need to use the GNU Make program, which is on the Skunkware CD. By default, it installs as /usr/local/bin/make.

As of UnixWare 7.1.3 and above, the GNU Make program is the OSTK portion of the UDK CD, and is in /usr/gnu/bin/gmake.

16.7.6.3. Readline

The Readline library is on the Skunkware CD. But it is not included on the UnixWare 7.1 Skunkware CD. If you have the UnixWare 7.0.0 or 7.0.1 Skunkware CDs, you can install it from there. Otherwise, try http://www.sco.com/skunkware/.

By default, Readline installs into /usr/local/lib and /usr/local/include. However, the PostgreSQL configure program will not find it there without help. If you installed Readline, then use the following options to configure:

```
./configure --with-libraries=/usr/local/lib --with-includes=/usr/local/include
```

16.7.6.4. Using the UDK on OpenServer

If you are using the new Universal Development Kit (UDK) compiler on OpenServer, you need to specify the locations of the UDK libraries:

```
./configure --with-libraries=/udk/usr/lib --with-includes=/udk/usr/include
```

Putting these together with the Readline options from above:

```
./configure --with-libraries="/udk/usr/lib /usr/local/lib" --with-includes="/udk/usr/include /usr/local/include"
```

16.7.6.5. Reading the PostgreSQL Man Pages

By default, the PostgreSQL man pages are installed into /usr/local/pgsql/share/man. By default, UnixWare does not look there for man pages. To be able to read them you need to modify the MANPATH variable in /etc/default/man, for example:

```
MANPATH=/usr/lib/scohelp/%L/man:/usr/dt/man:/usr/man:/usr/share/man:scohelp:/usr/local/man:/usr/local/pgsql/share/man
```

On OpenServer, some extra research needs to be invested to make the man pages usable, because the man system is a bit different from other platforms. Currently, PostgreSQL will not install them at all.
16.7.6.6. C99 Issues with the 7.1.1b Feature Supplement

For compilers earlier than the one released with OpenUNIX 8.0.0 (UnixWare 7.1.2), including the 7.1.1b Feature Supplement, you may need to specify `-Xb` in `CFLAGS` or the `CC` environment variable. The indication of this is an error in compiling `tuplesort.c` referencing inline functions. Apparently there was a change in the 7.1.2(8.0.0) compiler and beyond.

16.7.6.7. Threading on UnixWare

For threading, you must use `-Kpthread` on all `libpq`-using programs. `libpq` uses `pthread_*` calls, which are only available with the `-Kpthread/-Kthread` flag.

16.7.7. Solaris

PostgreSQL is well-supported on Solaris. The more up to date your operating system, the fewer issues you will experience; details below.

16.7.7.1. Required Tools

You can build with either GCC or Sun's compiler suite. For better code optimization, Sun's compiler is strongly recommended on the SPARC architecture. We have heard reports of problems when using GCC 2.95.1; GCC 2.95.3 or later is recommended. If you are using Sun's compiler, be careful not to select `/usr/ucb/cc`; use `/opt/SUNWspro/bin/cc`.

You can download Sun Studio from http://www.oracle.com/technetwork/server-storage/solarisstudio/downloads/. Many of GNU tools are integrated into Solaris 10, or they are present on the Solaris companion CD. If you like packages for older version of Solaris, you can find these tools at http://www.sunfreeware.com. If you prefer sources, look at http://www.gnu.org/order/ftp.html.

16.7.7.2. Problems with OpenSSL

When you build PostgreSQL with OpenSSL support you might get compilation errors in the following files:

- `src/backend/libpq/crypt.c`
- `src/backend/libpq/password.c`
- `src/interfaces/libpq/fe-auth.c`
- `src/interfaces/libpq/fe-connect.c`

This is because of a namespace conflict between the standard `/usr/include/crypt.h` header and the header files provided by OpenSSL.

Upgrading your OpenSSL installation to version 0.9.6a fixes this problem. Solaris 9 and above has a newer version of OpenSSL.
16.7.7.3. configure Complains About a Failed Test Program

If configure complains about a failed test program, this is probably a case of the run-time linker being unable to find some library, probably libz, libreadline or some other non-standard library such as libssl. To point it to the right location, set the LDFLAGS environment variable on the configure command line, e.g.,

```
configure ... LDFLAGS="-R /usr/sfw/lib:/opt/sfw/lib:/usr/local/lib"
```

See the ld man page for more information.

16.7.7.4. 64-bit Build Sometimes Crashes

On Solaris 7 and older, the 64-bit version of libc has a buggy vsnprintf routine, which leads to erratic core dumps in PostgreSQL. The simplest known workaround is to force PostgreSQL to use its own version of vsnprintf rather than the library copy. To do this, after you run configure edit a file produced by configure: In src/Makefile.global, change the line

```
LIBOBSJS =
```

to read

```
LIBOBSJS = snprintf.o
```

(There might be other files already listed in this variable. Order does not matter.) Then build as usual.

16.7.7.5. Compiling for Optimal Performance

On the SPARC architecture, Sun Studio is strongly recommended for compilation. Try using the -xO5 optimization flag to generate significantly faster binaries. Do not use any flags that modify behavior of floating-point operations and errno processing (e.g., -fast). These flags could raise some nonstandard PostgreSQL behavior for example in the date/time computing.

If you do not have a reason to use 64-bit binaries on SPARC, prefer the 32-bit version. The 64-bit operations are slower and 64-bit binaries are slower than the 32-bit variants. And on other hand, 32-bit code on the AMD64 CPU family is not native, and that is why 32-bit code is significant slower on this CPU family.

16.7.7.6. Using DTrace for Tracing PostgreSQL

Yes, using DTrace is possible. See Section 28.5 for further information. You can also find more information in this article: https://blogs.oracle.com/robertlor/entry/user_level_dtrace_probes_in.

If you see the linking of the postgres executable abort with an error message like:

```
Undefined first referenced
symbol in file
AbortTransaction utilities/probes.o
CommitTransaction utilities/probes.o
ld: fatal: Symbol referencing errors. No output written to postgres
collect2: ld returned 1 exit status
make: *** [postgres] Error 1
```
your DTrace installation is too old to handle probes in static functions. You need Solaris 10u4 or newer.
Chapter 17. Installation from Source Code on Windows

It is recommended that most users download the binary distribution for Windows, available as a graphical installer package from the PostgreSQL website. Building from source is only intended for people developing PostgreSQL or extensions.

There are several different ways of building PostgreSQL on Windows. The simplest way to build with Microsoft tools is to install Visual Studio 2019 and use the included compiler. It is also possible to build with the full Microsoft Visual C++ 2005 to 2019. In some cases that requires the installation of the Windows SDK in addition to the compiler.

It is also possible to build PostgreSQL using the GNU compiler tools provided by MinGW, or using Cygwin for older versions of Windows.

Finally, the client access library (libpq) can be built using Visual C++ 7.1 or Borland C++ for compatibility with statically linked applications built using these tools.

Building using MinGW or Cygwin uses the normal build system, see Chapter 16 and the specific notes in Section 16.7.5 and Section 16.7.2. To produce native 64 bit binaries in these environments, use the tools from MinGW-w64. These tools can also be used to cross-compile for 32 bit and 64 bit Windows targets on other hosts, such as Linux and Darwin. Cygwin is not recommended for running a production server, and it should only be used for running on older versions of Windows where the native build does not work, such as Windows 98. The official binaries are built using Visual Studio.

Native builds of psql don’t support command line editing. The Cygwin build does support command line editing, so it should be used where psql is needed for interactive use on Windows.

17.1. Building with Visual C++ or the Microsoft Windows SDK

PostgreSQL can be built using the Visual C++ compiler suite from Microsoft. These compilers can be either from Visual Studio, Visual Studio Express or some versions of the Microsoft Windows SDK. If you do not already have a Visual Studio environment set up, the easiest ways are to use the compilers from Visual Studio 2019 or those in the Windows SDK 10, which are both free downloads from Microsoft.


The tools for building using Visual C++ or Platform SDK are in the src/tools/msvc directory. When building, make sure there are no tools from MinGW or Cygwin present in your system PATH. Also, make sure you have all the required Visual C++ tools available in the PATH. In Visual Studio, start the Visual Studio Command Prompt. If you wish to build a 64-bit version, you must use the 64-bit version of the command, and vice versa. In the Microsoft Windows SDK, start the CMD shell listed under the SDK on the Start Menu. In recent SDK versions you can change the targeted CPU.
architecture, build type, and target OS by using the `setenv` command, e.g. `setenv /x86 /release /xp` to target Windows XP or later with a 32-bit release build. See `/?` for other options to `setenv`. All commands should be run from the `src\tools\msvc` directory.

Before you build, you may need to edit the file `config.pl` to reflect any configuration options you want to change, or the paths to any third party libraries to use. The complete configuration is determined by first reading and parsing the file `config_default.pl`, and then apply any changes from `config.pl`. For example, to specify the location of your Python installation, put the following in `config.pl`:

```
$config->{python} = 'c:\python26';
```

You only need to specify those parameters that are different from what’s in `config_default.pl`. If you need to set any other environment variables, create a file called `buildenv.pl` and put the required commands there. For example, to add the path for bison when it’s not in the PATH, create a file containing:

```
$ENV{PATH}=$ENV{PATH} . ';c:\some\where\bison\bin';
```

### 17.1.1. Requirements

The following additional products are required to build PostgreSQL. Use the `config.pl` file to specify which directories the libraries are available in.

**Microsoft Windows SDK**

If your build environment doesn’t ship with a supported version of the Microsoft Windows SDK it is recommended that you upgrade to the latest version (currently version 10), available for download from [https://www.microsoft.com/download](https://www.microsoft.com/download).

You must always include the Windows Headers and Libraries part of the SDK. If you install a Windows SDK including the Visual C++ Compilers, you don’t need Visual Studio to build. Note that as of Version 8.0a the Windows SDK no longer ships with a complete command-line build environment.

**ActiveState Perl**

ActiveState Perl is required to run the build generation scripts. MinGW or Cygwin Perl will not work. It must also be present in the PATH. Binaries can be downloaded from [http://www.activestate.com](http://www.activestate.com) (Note: version 5.8.3 or later is required, the free Standard Distribution is sufficient).

The following additional products are not required to get started, but are required to build the complete package. Use the `config.pl` file to specify which directories the libraries are available in.

**ActiveState TCL**

Required for building PL/TCL (Note: version 8.4 is required, the free Standard Distribution is sufficient).

**Bison and Flex**

Bison and Flex are required to build from Git, but not required when building from a release file. Only Bison 1.875 or versions 2.2 and later will work. Flex must be version 2.5.31 or later.
Both Bison and Flex are included in the msys tool suite, available from http://www.mingw.org/wiki/MSYS as part of the MinGW compiler suite.

You will need to add the directory containing flex.exe and bison.exe to the PATH environment variable in buildenv.pl unless they are already in PATH. In the case of MinGW, the directory is the \msys\1.0\bin subdirectory of your MinGW installation directory.

**Note:** The Bison distribution from GnuWin32 appears to have a bug that causes Bison to malfunction when installed in a directory with spaces in the name, such as the default location on English installations C:\Program Files\GnuWin32. Consider installing into C:\GnuWin32 or use the NTFS short name path to GnuWin32 in your PATH environment setting (e.g. C:\PROGRA~1\GnuWin32).

**Note:** The obsolete winflex binaries distributed on the PostgreSQL FTP site and referenced in older documentation will fail with “flex: fatal internal error, exec failed” on 64-bit Windows hosts. Use Flex from MSYS instead.

**Diff**

Diff is required to run the regression tests, and can be downloaded from http://gnuwin32.sourceforge.net.

**Gettext**

Gettext is required to build with NLS support, and can be downloaded from http://gnuwin32.sourceforge.net. Note that binaries, dependencies and developer files are all needed.

**MIT Kerberos**

Required for GSSAPI authentication support. MIT Kerberos can be downloaded from http://web.mit.edu/Kerberos/dist/index.html.

**libxml2 and libxslt**

Required for XML support. Binaries can be downloaded from http://zlatkovic.com/pub/libxml or source from http://xmlsoft.org. Note that libxml2 requires iconv, which is available from the same download location.

**openssl**


**ossu-uuid**

Required for UUID-OSSP support (contrib only). Source can be downloaded from http://www.ossu.org/pkg/lib/uuid/.

**Python**


**zlib**

Required for compression support in pg_dump and pg_restore. Binaries can be downloaded from http://www.zlib.net.
17.1.2. Special Considerations for 64-bit Windows

PostgreSQL will only build for the x64 architecture on 64-bit Windows, there is no support for Itanium processors.

Mixing 32- and 64-bit versions in the same build tree is not supported. The build system will automatically detect if it's running in a 32- or 64-bit environment, and build PostgreSQL accordingly. For this reason, it is important to start the correct command prompt before building.

To use a server-side third party library such as python or openssl, this library must also be 64-bit. There is no support for loading a 32-bit library in a 64-bit server. Several of the third party libraries that PostgreSQL supports may only be available in 32-bit versions, in which case they cannot be used with 64-bit PostgreSQL.

17.1.3. Building

To build all of PostgreSQL in release configuration (the default), run the command:

```bash
build
```

To build all of PostgreSQL in debug configuration, run the command:

```bash
build DEBUG
```

To build just a single project, for example psql, run the commands:

```bash
build psql
build DEBUG psql
```

To change the default build configuration to debug, put the following in the `buildenv.pl` file:

```bash
$ENV{CONFIG}="Debug";
```

It is also possible to build from inside the Visual Studio GUI. In this case, you need to run:

```bash
perl mkvcbuild.pl
```

from the command prompt, and then open the generated `pgsql.sln` (in the root directory of the source tree) in Visual Studio.

17.1.4. Cleaning and Installing

Most of the time, the automatic dependency tracking in Visual Studio will handle changed files. But if there have been large changes, you may need to clean the installation. To do this, simply run the `clean.bat` command, which will automatically clean out all generated files. You can also run it with the `dist` parameter, in which case it will behave like `make distclean` and remove the flex/bison output files as well.

By default, all files are written into a subdirectory of the `debug` or `release` directories. To install these files using the standard layout, and also generate the files required to initialize and use the database, run the command:

```bash
install c:\destination\directory
```
If you want to install only the client applications and interface libraries, then you can use these commands:

```
install c:\destination\directory client
```

### 17.1.5. Running the Regression Tests

To run the regression tests, make sure you have completed the build of all required parts first. Also, make sure that the DLLs required to load all parts of the system (such as the Perl and Python DLLs for the procedural languages) are present in the system path. If they are not, set it through the `buildenv.pl` file. To run the tests, run one of the following commands from the `src\tools\msvc` directory:

```
vcregress check
vcregress installcheck
vcregress plcheck
vcregress contribcheck
vcregress modulescheck
vcregress ecpgcheck
vcregress isolationcheck
vcregress bincheck
vcregress recoverycheck
vcregress upgradecheck
```

To change the schedule used (default is parallel), append it to the command line like:

```
vcregress check serial
```

For more information about the regression tests, see Chapter 31.

Running the regression tests on client programs, with `vcregress bincheck`, or on recovery tests, with `vcregress recoverycheck`, requires an additional Perl module to be installed:

**IPC::Run**

As of this writing, **IPC::Run** is not included in the ActiveState Perl installation, nor in the ActiveState Perl Package Manager (PPM) library. To install, download the `IPC-Run-<version>.tar.gz` source archive from CPAN, at https://metacpan.org/release/IPC-Run/, and uncompress. Edit the `buildenv.pl` file, and add a `PERL5LIB` variable to point to the `lib` subdirectory from the extracted archive. For example:

```
$ENV{PERL5LIB}=$ENV{PERL5LIB} . ';c:\IPC-Run-0.94\lib';
```
17.1.6. Building the Documentation

Building the PostgreSQL documentation in HTML format requires several tools and files. Create a root directory for all these files, and store them in the subdirectories in the list below.

OpenJade 1.3.1-2

- Download from http://sourceforge.net/projects/openjade/files/openjade/1.3.1/openjade-1_3_1-2-bin.zip/download and uncompress in the subdirectory openjade-1.3.1.

DocBook DTD 4.2


DocBook DSSSL 1.79


ISO character entities


Edit the buildenv.pl file, and add a variable for the location of the root directory, for example:

```bash
$ENV{DOCROOT}='c:\docbook';
```

To build the documentation, run the command `builddoc.bat`. Note that this will actually run the build twice, in order to generate the indexes. The generated HTML files will be in `doc\src\sgml`.

17.2. Building libpq with Visual C++ or Borland C++

Using Visual C++ 7.1-9.0 or Borland C++ to build libpq is only recommended if you need a version with different debug/release flags, or if you need a static library to link into an application. For normal use the MinGW or Visual Studio or Windows SDK method is recommended.

To build the libpq client library using Visual Studio 7.1 or later, change into the `src` directory and type the command:

```bash
nmake /f win32.mak
```

To build a 64-bit version of the libpq client library using Visual Studio 8.0 or later, change into the `src` directory and type in the command:

```bash
nmake /f win32.mak CPU=AMD64
```

See the `win32.mak` file for further details about supported variables.

To build the libpq client library using Borland C++, change into the `src` directory and type the command:

```bash
make -N -DCFG=Release /f bcc32.mak
```
17.2.1. Generated Files

The following files will be built:

interfaces\libpq\Release\libpq.dll

The dynamically linkable frontend library

interfaces\libpq\Release\libpqdll.lib

Import library to link your programs to libpq.dll

interfaces\libpq\Release\libpq.lib

Static version of the frontend library

Normally you do not need to install any of the client files. You should place the libpq.dll file in the same directory as your applications executable file. Do not install libpq.dll into your Windows, System or System32 directory unless absolutely necessary. If this file is installed using a setup program, then it should be installed with version checking using the VERSIONINFO resource included in the file, to ensure that a newer version of the library is not overwritten.

If you are planning to do development using libpq on this machine, you will have to add the src\include and src\interfaces\libpq subdirectories of the source tree to the include path in your compiler’s settings.

To use the library, you must add the libpqdll.lib file to your project. (In Visual C++, just right-click on the project and choose to add it.)
Chapter 18. Server Setup and Operation

This chapter discusses how to set up and run the database server and its interactions with the operating system.

18.1. The PostgreSQL User Account

As with any server daemon that is accessible to the outside world, it is advisable to run PostgreSQL under a separate user account. This user account should only own the data that is managed by the server, and should not be shared with other daemons. (For example, using the user `nobody` is a bad idea.) It is not advisable to install executables owned by this user because compromised systems could then modify their own binaries.

To add a Unix user account to your system, look for a command `useradd` or `adduser`. The user name `postgres` is often used, and is assumed throughout this book, but you can use another name if you like.

18.2. Creating a Database Cluster

Before you can do anything, you must initialize a database storage area on disk. We call this a database cluster. (The SQL standard uses the term catalog cluster.) A database cluster is a collection of databases that is managed by a single instance of a running database server. After initialization, a database cluster will contain a database named `postgres`, which is meant as a default database for use by utilities, users and third party applications. The database server itself does not require the `postgres` database to exist, but many external utility programs assume it exists. Another database created within each cluster during initialization is called `template1`. As the name suggests, this will be used as a template for subsequently created databases; it should not be used for actual work. (See Chapter 22 for information about creating new databases within a cluster.)

In file system terms, a database cluster is a single directory under which all data will be stored. We call this the data directory or data area. It is completely up to you where you choose to store your data. There is no default, although locations such as `/usr/local/pgsql/data` or `/var/lib/pgsql/data` are popular. To initialize a database cluster, use the command `initdb`, which is installed with PostgreSQL. The desired file system location of your database cluster is indicated by the `-D` option, for example:

```
$ initdb -D /usr/local/pgsql/data
```

Note that you must execute this command while logged into the PostgreSQL user account, which is described in the previous section.

**Tip:** As an alternative to the `-D` option, you can set the environment variable `PGDATA`.

Alternatively, you can run `initdb` via the `pg_ctl` program like so:

```
$ pg_ctl -D /usr/local/pgsql/data initdb
```

This may be more intuitive if you are using `pg_ctl` for starting and stopping the server (see Section 18.3), so that `pg_ctl` would be the sole command you use for managing the database server instance.
Chapter 18. Server Setup and Operation

initdb will attempt to create the directory you specify if it does not already exist. Of course, this will fail if initdb does not have permissions to write in the parent directory. It’s generally recommendable that the PostgreSQL user own not just the data directory but its parent directory as well, so that this should not be a problem. If the desired parent directory doesn’t exist either, you will need to create it first, using root privileges if the grandparent directory isn’t writable. So the process might look like this:

root# mkdir /usr/local/pgsql
root# chown postgres /usr/local/pgsql
root# su postgres
postgres$ initdb -D /usr/local/pgsql/data

initdb will refuse to run if the data directory exists and already contains files; this is to prevent accidentally overwriting an existing installation.

Because the data directory contains all the data stored in the database, it is essential that it be secured from unauthorized access. initdb therefore revokes access permissions from everyone but the PostgreSQL user.

However, while the directory contents are secure, the default client authentication setup allows any local user to connect to the database and even become the database superuser. If you do not trust other local users, we recommend you use one of initdb’s -W, --pwprompt or --pwfile options to assign a password to the database superuser. Also, specify -A md5 or --A password so that the default trust authentication mode is not used; or modify the generated pg_hba.conf file after running initdb, but before you start the server for the first time. (Other reasonable approaches include using peer authentication or file system permissions to restrict connections. See Chapter 20 for more information.)

initdb also initializes the default locale for the database cluster. Normally, it will just take the locale settings in the environment and apply them to the initialized database. It is possible to specify a different locale for the database; more information about that can be found in Section 23.1. The default sort order used within the particular database cluster is set by initdb, and while you can create new databases using different sort order, the order used in the template databases that initdb creates cannot be changed without dropping and recreating them. There is also a performance impact for using locales other than C or POSIX. Therefore, it is important to make this choice correctly the first time.

initdb also sets the default character set encoding for the database cluster. Normally this should be chosen to match the locale setting. For details see Section 23.3.

Non-C and non-POSIX locales rely on the operating system’s collation library for character set ordering. This controls the ordering of keys stored in indexes. For this reason, a cluster cannot switch to an incompatible collation library version, either through snapshot restore, binary streaming replication, a different operating system, or an operating system upgrade.

18.2.1. Use of Secondary File Systems

Many installations create their database clusters on file systems (volumes) other than the machine’s “root” volume. If you choose to do this, it is not advisable to try to use the secondary volume’s topmost directory (mount point) as the data directory. Best practice is to create a directory within the mount-point directory that is owned by the PostgreSQL user, and then create the data directory within that. This avoids permissions problems, particularly for operations such as pg_upgrade, and it also ensures clean failures if the secondary volume is taken offline.
18.2.2. Use of Network File Systems

Many installations create their database clusters on network file systems. Sometimes this is done via NFS, or by using a Network Attached Storage (NAS) device that uses NFS internally. PostgreSQL does nothing special for NFS file systems, meaning it assumes NFS behaves exactly like locally-connected drives. If the client or server NFS implementation does not provide standard file system semantics, this can cause reliability problems (see http://www.time-travellers.org/shane/papers/NFS_considered_harmful.html). Specifically, delayed (asynchronous) writes to the NFS server can cause data corruption problems. If possible, mount the NFS file system synchronously (without caching) to avoid this hazard. Also, soft-mounting the NFS file system is not recommended.

Storage Area Networks (SAN) typically use communication protocols other than NFS, and may or may not be subject to hazards of this sort. It’s advisable to consult the vendor’s documentation concerning data consistency guarantees. PostgreSQL cannot be more reliable than the file system it’s using.

18.3. Starting the Database Server

Before anyone can access the database, you must start the database server. The database server program is called postgres. The postgres program must know where to find the data it is supposed to use. This is done with the \texttt{-D} option. Thus, the simplest way to start the server is:

\begin{verbatim}
$ postgres -D /usr/local/pgsql/data
\end{verbatim}

which will leave the server running in the foreground. This must be done while logged into the PostgreSQL user account. Without \texttt{-D}, the server will try to use the data directory named by the environment variable \texttt{PGDATA}. If that variable is not provided either, it will fail.

Normally it is better to start postgres in the background. For this, use the usual Unix shell syntax:

\begin{verbatim}
$ postgres -D /usr/local/pgsql/data >logfile 2>&1 &
\end{verbatim}

It is important to store the server’s stdout and stderr output somewhere, as shown above. It will help for auditing purposes and to diagnose problems. (See Section 24.3 for a more thorough discussion of log file handling.)

The postgres program also takes a number of other command-line options. For more information, see the postgres reference page and Chapter 19 below.

This shell syntax can get tedious quickly. Therefore the wrapper program \texttt{pg_ctl} is provided to simplify some tasks. For example:

\begin{verbatim}
png_ctl start -l logfile
\end{verbatim}

will start the server in the background and put the output into the named log file. The \texttt{-D} option has the same meaning here as for postgres. \texttt{pg_ctl} is also capable of stopping the server.

Normally, you will want to start the database server when the computer boots. Autostart scripts are operating-system-specific. There are a few distributed with PostgreSQL in the contrib/start-scripts directory. Installing one will require root privileges.

Different systems have different conventions for starting up daemons at boot time. Many systems have a file /etc/rc.local or /etc/rc.d/rc.local. Others use init.d or rc.d directories. Whatever you do, the server must be run by the PostgreSQL user account \textit{and not by root} or any
other user. Therefore you probably should form your commands using `su postgres -c '...'
For example:

```
su postgres -c 'pg_ctl start -D /usr/local/pgsql/data -l serverlog'
```

Here are a few more operating-system-specific suggestions. (In each case be sure to use the proper
installation directory and user name where we show generic values.)

- For FreeBSD, look at the file `contrib/start-scripts/freebsd` in the PostgreSQL source
distribution.
- On OpenBSD, add the following lines to the file `/etc/rc.local`

  ```
  if [ -x /usr/local/pgsql/bin/pg_ctl -a -x /usr/local/pgsql/bin/postgres ]; then
    su -l postgres -c '/usr/local/pgsql/bin/pg_ctl start -s -l /var/postgresql/log -D /
    echo -n ' postgresql'
  fi
  ```
- On Linux systems either add `/usr/local/pgsql/bin/pg_ctl start -l logfile -D /usr/local/pgsql/data`
to `/etc/rc.d/rc.local` or `/etc/rc.local` or look at the file `contrib/start-scripts/linux` in the PostgreSQL source distribution.

When using systemd, you can use the following service unit file (e.g., at
`/etc/systemd/system/postgresql.service`):

```
[Unit]
Description=PostgreSQL database server
Documentation=man:postgres(1)

[Service]
Type=notify
User=postgres
ExecStart=/usr/local/pgsql/bin/postgres -D /usr/local/pgsql/data
ExecReload=/bin/kill -HUP $MAINPID
KillMode=mixed
KillSignal=SIGINT
TimeoutSec=0

[Install]
WantedBy=multi-user.target
```

Using `Type=notify` requires that the server binary was built with `configure
--with-systemd`.

Consider carefully the timeout setting. systemd has a default timeout of 90 seconds as of this writing
and will kill a process that does not notify readiness within that time. But a PostgreSQL server that
might have to perform crash recovery at startup could take much longer to become ready. The
suggested value of 0 disables the timeout logic.

- On NetBSD, use either the FreeBSD or Linux start scripts, depending on preference.
- On Solaris, create a file called `/etc/init.d/postgresql` that contains the following line:

  ```
su - postgres -c "*/usr/local/pgsql/bin/pg_ctl start -l logfile -D /usr/local/pgsql/data"
  ```

Then, create a symbolic link to it in `/etc/rc3.d` as S99postgresql.
While the server is running, its PID is stored in the file `postmaster.pid` in the data directory. This is used to prevent multiple server instances from running in the same data directory and can also be used for shutting down the server.

### 18.3.1. Server Start-up Failures

There are several common reasons the server might fail to start. Check the server’s log file, or start it by hand (without redirecting standard output or standard error) and see what error messages appear. Below we explain some of the most common error messages in more detail.

```
LOG: could not bind IPv4 socket: Address already in use
HINT: Is another postmaster already running on port 5432? If not, wait a few seconds and retry.
FATAL: could not create TCP/IP listen socket
```

This usually means just what it suggests: you tried to start another server on the same port where one is already running. However, if the kernel error message is not `Address already in use` or some variant of that, there might be a different problem. For example, trying to start a server on a reserved port number might draw something like:

```
$ postgres -p 666
LOG: could not bind IPv4 socket: Permission denied
HINT: Is another postmaster already running on port 666? If not, wait a few seconds and retry.
FATAL: could not create TCP/IP listen socket
```

A message like:

```
FATAL: could not create shared memory segment: Invalid argument
DETAIL: Failed system call was shmget(key=5440001, size=4011376640, 03600).
```

probably means your kernel’s limit on the size of shared memory is smaller than the work area PostgreSQL is trying to create (4011376640 bytes in this example). Or it could mean that you do not have System-V-style shared memory support configured into your kernel at all. As a temporary workaround, you can try starting the server with a smaller-than-normal number of buffers (shared_buffers). You will eventually want to reconfigure your kernel to increase the allowed shared memory size. You might also see this message when trying to start multiple servers on the same machine, if their total space requested exceeds the kernel limit.

An error like:

```
FATAL: could not create semaphores: No space left on device
DETAIL: Failed system call was semget(5440126, 17, 03600).
```

does not mean you’ve run out of disk space. It means your kernel’s limit on the number of System V semaphores is smaller than the number PostgreSQL wants to create. As above, you might be able to work around the problem by starting the server with a reduced number of allowed connections (max_connections), but you’ll eventually want to increase the kernel limit.

If you get an “illegal system call” error, it is likely that shared memory or semaphores are not supported in your kernel at all. In that case your only option is to reconfigure the kernel to enable these features.

Details about configuring System V IPC facilities are given in Section 18.4.1.
18.3.2. Client Connection Problems

Although the error conditions possible on the client side are quite varied and application-dependent, a few of them might be directly related to how the server was started. Conditions other than those shown below should be documented with the respective client application.

psql: could not connect to server: Connection refused
Is the server running on host "server.joe.com" and accepting TCP/IP connections on port 5432?

This is the generic “I couldn’t find a server to talk to” failure. It looks like the above when TCP/IP communication is attempted. A common mistake is to forget to configure the server to allow TCP/IP connections.

Alternatively, you’ll get this when attempting Unix-domain socket communication to a local server:

psql: could not connect to server: No such file or directory
Is the server running locally and accepting connections on Unix domain socket "/tmp/.s.PGSQL.5432"?

The last line is useful in verifying that the client is trying to connect to the right place. If there is in fact no server running there, the kernel error message will typically be either Connection refused or No such file or directory, as illustrated. (It is important to realize that Connection refused in this context does not mean that the server got your connection request and rejected it. That case will produce a different message, as shown in Section 20.4.) Other error messages such as Connection timed out might indicate more fundamental problems, like lack of network connectivity.

18.4. Managing Kernel Resources

PostgreSQL can sometimes exhaust various operating system resource limits, especially when multiple copies of the server are running on the same system, or in very large installations. This section explains the kernel resources used by PostgreSQL and the steps you can take to resolve problems related to kernel resource consumption.

18.4.1. Shared Memory and Semaphores

Shared memory and semaphores are collectively referred to as “System V IPC” (together with message queues, which are not relevant for PostgreSQL). Except on Windows, where PostgreSQL provides its own replacement implementation of these facilities, these facilities are required in order to run PostgreSQL.

The complete lack of these facilities is usually manifested by an Illegal system call error upon server start. In that case there is no alternative but to reconfigure your kernel. PostgreSQL won’t work without them. This situation is rare, however, among modern operating systems.

When PostgreSQL exceeds one of the various hard IPC limits, the server will refuse to start and should leave an instructive error message describing the problem and what to do about it. (See also Section 18.3.1.) The relevant kernel parameters are named consistently across different systems; Table 18-1 gives an overview. The methods to set them, however, vary. Suggestions for some platforms are given below.
**Chapter 18. Server Setup and Operation**

**Note:** Prior to PostgreSQL 9.3, the amount of System V shared memory required to start the server was much larger. If you are running an older version of the server, please consult the documentation for your server version.

### Table 18-1. System V IPC Parameters

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
<th>Reasonable values</th>
</tr>
</thead>
<tbody>
<tr>
<td>SHMMAX</td>
<td>Maximum size of shared memory segment (bytes)</td>
<td>at least 1kB (more if running many copies of the server)</td>
</tr>
<tr>
<td>SHMMIN</td>
<td>Minimum size of shared memory segment (bytes)</td>
<td>1</td>
</tr>
<tr>
<td>SHMALL</td>
<td>Total amount of shared memory available (bytes or pages)</td>
<td>if bytes, same as SHMMAX; if pages, ceil(SHMMAX/PAGE_SIZE)</td>
</tr>
<tr>
<td>SHMSEG</td>
<td>Maximum number of shared memory segments per process</td>
<td>only 1 segment is needed, but the default is much higher</td>
</tr>
<tr>
<td>SHMMNI</td>
<td>Maximum number of shared memory segments system-wide</td>
<td>like SHMSEG plus room for other applications</td>
</tr>
<tr>
<td>SEMMNI</td>
<td>Maximum number of semaphore identifiers (i.e., sets)</td>
<td>at least ceil((max_connections + autovacuum_max_workers + max_worker_processes + 5) / 16)</td>
</tr>
<tr>
<td>SEMMNS</td>
<td>Maximum number of semaphores system-wide</td>
<td>ceil((max_connections + autovacuum_max_workers + max_worker_processes + 5) / 16) * 17 plus room for other applications</td>
</tr>
<tr>
<td>SEMMSL</td>
<td>Maximum number of semaphores per set</td>
<td>at least 17</td>
</tr>
<tr>
<td>SEMMAP</td>
<td>Number of entries in semaphore map</td>
<td>see text</td>
</tr>
<tr>
<td>SEMVMX</td>
<td>Maximum value of semaphore</td>
<td>at least 1000 (The default is often 32767; do not change unless necessary)</td>
</tr>
</tbody>
</table>

PostgreSQL requires a few bytes of System V shared memory (typically 48 bytes, on 64-bit platforms) for each copy of the server. On most modern operating systems, this amount can easily be allocated. However, if you are running many copies of the server, or if other applications are also using System V shared memory, it may be necessary to increase SHMMAX, the maximum size in bytes of a shared memory segment, or SHMALL, the total amount of System V shared memory system-wide. Note that SHMALL is measured in pages rather than bytes on many systems.

Less likely to cause problems is the minimum size for shared memory segments (SHMMIN), which should be at most approximately 32 bytes for PostgreSQL (it is usually just 1). The maximum number of segments system-wide (SHMMNI) or per-process (SHMSEG) are unlikely to cause a problem unless your system has them set to zero.

PostgreSQL uses one semaphore per allowed connection (max_connections), allowed
autovacuum worker process (autovacuum_max_workers) and allowed background process
(max_worker_processes), in sets of 16. Each such set will also contain a 17th semaphore
which contains a “magic number”, to detect collision with semaphore sets used by other
applications. The maximum number of semaphores in the system is set by SEMMNS, which
consequently must be at least as high as max_connections plus autovacuum_max_workers
plus max_worker_processes, plus one extra for each 16 allowed connections plus workers
(see the formula in Table 18-1). The parameter SEMMNI determines the limit on the number of
semaphore sets that can exist on the system at one time. Hence this parameter must be at least
\[
\text{ceil}(\text{max_connections} + \text{autovacuum_max_workers} + \text{max_worker_processes} + 5) / 16
\]
Lowering the number of allowed connections is a temporary workaround for failures, which
are usually confusingly worded “No space left on device”, from the function semget.
In some cases it might also be necessary to increase SEMMAP to be at least on the order of SEMMNS.
If the system has this parameter (many do not), it defines the size of the semaphore resource map,
in which each contiguous block of available semaphores needs an entry. When a semaphore set is
freed it is either added to an existing entry that is adjacent to the freed block or it is registered under
a new map entry. If the map is full, the freed semaphores get lost (until reboot). Fragmentation of the
semaphore space could over time lead to fewer available semaphores than there should be.
The SEMMSL parameter, which determines how many semaphores can be in a set, must be at least 17
for PostgreSQL.
Various other settings related to “semaphore undo”, such as SEMMNU and SEMUME, do not affect Post-
geSQL.
AIX
At least as of version 5.1, it should not be necessary to do any special configuration for such
parameters as SHMMAX, as it appears this is configured to allow all memory to be used as shared
memory. That is the sort of configuration commonly used for other databases such as DB/2.
It might, however, be necessary to modify the global ulimit information in
/etc/security/limits, as the default hard limits for file sizes (fsize) and numbers of files
(nofiles) might be too low.
FreeBSD
The default IPC settings can be changed using the sysctl or loader interfaces. The following
parameters can be set using sysctl:

```
# sysctl kern.ipc.shmall=32768
# sysctl kern.ipc.shmmax=134217728
```
To make these settings persist over reboots, modify /etc/sysctl.conf.
These semaphore-related settings are read-only as far as sysctl is concerned, but can be set in
/boot/loader.conf:
kern.ipc.semnni=256
kern.ipc.semms=512
After modifying that file, a reboot is required for the new settings to take effect.
You might also want to configure your kernel to lock shared memory into RAM and prevent
it from being paged out to swap. This can be accomplished using the sysctl setting
kern.ipc.shm_use_phys.
If running in FreeBSD jails by enabling sysctl’s security.jail.sysv_ipc_allowed, post-
masters running in different jails should be run by different operating system users. This
improves security because it prevents non-root users from interfering with shared memory or
semaphores in different jails, and it allows the PostgreSQL IPC cleanup code to function prop-
erly. (In FreeBSD 6.0 and later the IPC cleanup code does not properly detect processes in other jails, preventing the running of postmasters on the same port in different jails.)

FreeBSD versions before 4.0 work like old OpenBSD (see below).

NetBSD

In NetBSD 5.0 and later, IPC parameters can be adjusted using `sysctl`, for example:

```bash
# sysctl -w kern.ipc.semmni=100
```

To make these settings persist over reboots, modify `/etc/sysctl.conf`.

You will usually want to increase `kern.ipc.semmni` and `kern.ipc.semmns`, as NetBSD’s default settings for these are uncomfortably small.

You might also want to configure your kernel to lock shared memory into RAM and prevent it from being paged out to swap. This can be accomplished using the `sysctl` setting `kern.ipc.shm_use_phys`.

NetBSD versions before 5.0 work like old OpenBSD (see below), except that kernel parameters should be set with the keyword `options` not `option`.

OpenBSD

In OpenBSD 3.3 and later, IPC parameters can be adjusted using `sysctl`, for example:

```bash
# sysctl kern.seminfo.semmni=100
```

To make these settings persist over reboots, modify `/etc/sysctl.conf`.

You will usually want to increase `kern.seminfo.semmni` and `kern.seminfo.semmns`, as OpenBSD’s default settings for these are uncomfortably small.

In older OpenBSD versions, you will need to build a custom kernel to change the IPC parameters. Make sure that the options `SYSVSHM` and `SYSVSEM` are enabled, too. (They are by default.) The following shows an example of how to set the various parameters in the kernel configuration file:

```plaintext
option     SYSVSHM
option     SHMMAXPGS=4096
option     SHMSEG=256

option     SYSVSEM
option     SEMMNI=256
option     SEMMNS=512
option     SEMMNU=256
```

HP-UX

The default settings tend to suffice for normal installations. On HP-UX 10, the factory default for `SEMMNS` is 128, which might be too low for larger database sites.

IPC parameters can be set in the System Administration Manager (SAM) under Kernel Configuration → Configurable Parameters. Choose Create A New Kernel when you’re done.

Linux

The default maximum segment size is 32 MB, and the default maximum total size is 2097152 pages. A page is almost always 4096 bytes except in unusual kernel configurations with “huge pages” (use `getconf PAGE_SIZE` to verify).

The shared memory size settings can be changed via the `sysctl` interface. For example, to allow 16 GB:

```bash
$ sysctl -w kernel.shmmax=17179869184
$ sysctl -w kernel.shmall=4194304
```
In addition these settings can be preserved between reboots in the file `/etc/sysctl.conf`. Doing that is highly recommended.

Ancient distributions might not have the `sysctl` program, but equivalent changes can be made by manipulating the `/proc` file system:

```
$ echo 17179869184 > /proc/sys/kernel/shmmax
$ echo 4194304 > /proc/sys/kernel/shmall
```

The remaining defaults are quite generously sized, and usually do not require changes.

**OS X**

The recommended method for configuring shared memory in OS X is to create a file named `/etc/sysctl.conf`, containing variable assignments such as:

```
kern.sysv.shmmax=4194304
kern.sysv.shmmin=1
kern.sysv.shmmni=32
kern.sysv.shmseg=8
kern.sysv.shmall=1024
```

Note that in some OS X versions, *all five* shared-memory parameters must be set in `/etc/sysctl.conf`, else the values will be ignored.

Beware that recent releases of OS X ignore attempts to set `SHMMAX` to a value that isn’t an exact multiple of 4096.

`SHMALL` is measured in 4 kB pages on this platform.

In older OS X versions, you will need to reboot to have changes in the shared memory parameters take effect. As of 10.5 it is possible to change all but `SHMMNI` on the fly, using `sysctl`. But it’s still best to set up your preferred values via `/etc/sysctl.conf`, so that the values will be kept across reboots.

The file `/etc/sysctl.conf` is only honored in OS X 10.3.9 and later. If you are running a previous 10.3.x release, you must edit the file `/etc/rc` and change the values in the following commands:

```
sysctl -w kern.sysv.shmmax
sysctl -w kern.sysv.shmmin
sysctl -w kern.sysv.shmmni
sysctl -w kern.sysv.shmseg
sysctl -w kern.sysv.shmall
```

Note that `/etc/rc` is usually overwritten by OS X system updates, so you should expect to have to redo these edits after each update.

In OS X 10.2 and earlier, instead edit these commands in the file `/System/Library/StartupItems/SystemTuning/SystemTuning`

**SCO OpenServer**

In the default configuration, only 512 kB of shared memory per segment is allowed. To increase the setting, first change to the directory `/etc/conf/cf.d`. To display the current value of `SHMMAX`, run:

```
./configure -y SHMMAX
```

To set a new value for `SHMMAX`, run:

```
./configure SHMMAX=value
```

where `value` is the new value you want to use (in bytes). After setting `SHMMAX`, rebuild the kernel:

```
./link_unix
```

and reboot.
Solaris 2.6 to 2.9 (Solaris 6 to Solaris 9)

The relevant settings can be changed in /etc/system, for example:

```plaintext
set shmsys:shminfo_shmmax=0x2000000
set shmsys:shminfo_shmmin=1
set shmsys:shminfo_shmmni=256
set shmsys:shminfo_shmseg=256

set semsys:seminfo_semmap=256
set semsys:seminfo_semmni=512
set semsys:seminfo_semmns=512
set semsys:seminfo_semmns=32
```

You need to reboot for the changes to take effect. See also http://sunsite.uakom.sk/sunworldonline/swol-09-1997/swol-09-insidesolaris.html for information on shared memory under older versions of Solaris.

Solaris 2.10 (Solaris 10) and later

OpenSolaris

In Solaris 10 and later, and OpenSolaris, the default shared memory and semaphore settings are good enough for most PostgreSQL applications. Solaris now defaults to a \texttt{SHMMAX} of one-quarter of system RAM. To further adjust this setting, use a project setting associated with the \texttt{postgres} user. For example, run the following as root:

```plaintext
projadd -c "PostgreSQL DB User" -K "project.max-shm-memory=(privileged,8GB,deny)" -U postgres -G postgres user.postgres
```

This command adds the \texttt{user.postgres} project and sets the shared memory maximum for the \texttt{postgres} user to 8GB, and takes effect the next time that user logs in, or when you restart PostgreSQL (not reload). The above assumes that PostgreSQL is run by the \texttt{postgres} user in the \texttt{postgres} group. No server reboot is required.

Other recommended kernel setting changes for database servers which will have a large number of connections are:

```plaintext
project.max-shm-ids=(priv,32768,deny)
project.max-sem-ids=(priv,4096,deny)
project.max-msg-ids=(priv,4096,deny)
```

Additionally, if you are running PostgreSQL inside a zone, you may need to raise the zone resource usage limits as well. See "Chapter 2: Projects and Tasks" in the \textit{System Administrator's Guide} for more information on projects and \texttt{prctl}.

UnixWare

On UnixWare 7, the maximum size for shared memory segments is 512 kB in the default configuration. To display the current value of \texttt{SHMMAX}, run:

```plaintext
/etc/conf/bin/idtune -g SHMMAX
```

which displays the current, default, minimum, and maximum values. To set a new value for \texttt{SHMMAX}, run:

```plaintext
/etc/conf/bin/idtune SHMMAX value
```

where \texttt{value} is the new value you want to use (in bytes). After setting \texttt{SHMMAX}, rebuild the kernel:

```plaintext
/etc/conf/bin/idbuild -B
```

and reboot.
18.4.2. systemd RemoveIPC

If systemd is in use, some care must be taken that IPC resources (shared memory and semaphores) are not prematurely removed by the operating system. This is especially of concern when installing PostgreSQL from source. Users of distribution packages of PostgreSQL are less likely to be affected, as the postgres user is then normally created as a system user.

The setting `RemoveIPC` in `logind.conf` controls whether IPC objects are removed when a user fully logs out. System users are exempt. This setting defaults to on in stock systemd, but some operating system distributions default it to off.

A typical observed effect when this setting is on is that the semaphore objects used by a PostgreSQL server are removed at apparently random times, leading to the server crashing with log messages like

```
LOG: semctl(1234567890, 0, IPC_RMID, ...) failed: Invalid argument
```

Different types of IPC objects (shared memory vs. semaphores, System V vs. POSIX) are treated slightly differently by systemd, so one might observe that some IPC resources are not removed in the same way as others. But it is not advisable to rely on these subtle differences.

A “user logging out” might happen as part of a maintenance job or manually when an administrator logs in as the postgres user or something similar, so it is hard to prevent in general.

What is a “system user” is determined at systemd compile time from the `SYS_UID_MAX` setting in `/etc/login.defs`.

Packaging and deployment scripts should be careful to create the postgres user as a system user by using `useradd -r`, `adduser --system`, or equivalent.

Alternatively, if the user account was created incorrectly or cannot be changed, it is recommended to set

```
RemoveIPC=no
```

in `/etc/systemd/logind.conf` or another appropriate configuration file.

<table>
<thead>
<tr>
<th>Caution</th>
</tr>
</thead>
<tbody>
<tr>
<td>At least one of these two things has to be ensured, or the PostgreSQL server will be very unreliable.</td>
</tr>
</tbody>
</table>

18.4.3. Resource Limits

Unix-like operating systems enforce various kinds of resource limits that might interfere with the operation of your PostgreSQL server. Of particular importance are limits on the number of processes per user, the number of open files per process, and the amount of memory available to each process. Each of these have a “hard” and a “soft” limit. The soft limit is what actually counts but it can be changed by the user up to the hard limit. The hard limit can only be changed by the root user. The system call `setrlimit` is responsible for setting these parameters. The shell’s built-in command `ulimit` (Bourne shells) or `limit` (csh) is used to control the resource limits from the command line.

On BSD-derived systems the file `/etc/login.conf` controls the various resource limits set during login. See the operating system documentation for details. The relevant parameters are `maxproc`, `openfiles`, and `datasize`. For example:

```
default:
...
```
Kernels can also have system-wide limits on some resources.

- On Linux /proc/sys/fs/file-max determines the maximum number of open files that the kernel will support. It can be changed by writing a different number into the file or by adding an assignment in /etc/sysctl.conf. The maximum limit of files per process is fixed at the time the kernel is compiled; see /usr/src/linux/Documentation/proc.txt for more information.

The PostgreSQL server uses one process per connection so you should provide for at least as many processes as allowed connections, in addition to what you need for the rest of your system. This is usually not a problem but if you run several servers on one machine things might get tight.

The factory default limit on open files is often set to “socially friendly” values that allow many users to coexist on a machine without using an inappropriate fraction of the system resources. If you run many servers on a machine this is perhaps what you want, but on dedicated servers you might want to raise this limit.

On the other side of the coin, some systems allow individual processes to open large numbers of files; if more than a few processes do so then the system-wide limit can easily be exceeded. If you find this happening, and you do not want to alter the system-wide limit, you can set PostgreSQL’s max_files_per_process configuration parameter to limit the consumption of open files.

### 18.4.4. Linux Memory Overcommit

In Linux 2.4 and later, the default virtual memory behavior is not optimal for PostgreSQL. Because of the way that the kernel implements memory overcommit, the kernel might terminate the PostgreSQL postmaster (the master server process) if the memory demands of either PostgreSQL or another process cause the system to run out of virtual memory.

If this happens, you will see a kernel message that looks like this (consult your system documentation and configuration on where to look for such a message):

```
Out of Memory: Killed process 12345 (postgres).
```

This indicates that the postgres process has been terminated due to memory pressure. Although existing database connections will continue to function normally, no new connections will be accepted. To recover, PostgreSQL will need to be restarted.

One way to avoid this problem is to run PostgreSQL on a machine where you can be sure that other processes will not run the machine out of memory. If memory is tight, increasing the swap space of the operating system can help avoid the problem, because the out-of-memory (OOM) killer is invoked only when physical memory and swap space are exhausted.

If PostgreSQL itself is the cause of the system running out of memory, you can avoid the problem by changing your configuration. In some cases, it may help to lower memory-related configuration parameters, particularly shared_buffers and work_mem. In other cases, the problem may be caused by allowing too many connections to the database server itself. In many cases, it may be better to reduce max_connections and instead make use of external connection-pooling software.
Chapter 18. Server Setup and Operation

On Linux 2.6 and later, it is possible to modify the kernel’s behavior so that it will not “overcommit” memory. Although this setting will not prevent the OOM killer\(^1\) from being invoked altogether, it will lower the chances significantly and will therefore lead to more robust system behavior. This is done by selecting strict overcommit mode via `sysctl`:

```
sysctl -w vm.overcommit_memory=2
```

or placing an equivalent entry in `/etc/sysctl.conf`. You might also wish to modify the related setting `vm.overcommit_ratio`. For details see the kernel documentation file https://www.kernel.org/doc/Documentation/vm/overcommit-accounting.

Another approach, which can be used with or without altering `vm.overcommit_memory`, is to set the process-specific **OOM score adjustment** value for the postmaster process to `-1000`, thereby guaranteeing it will not be targeted by the OOM killer. The simplest way to do this is to execute

```
echo -1000 > /proc/self/oom_score_adj
```

in the postmaster’s startup script just before invoking the postmaster. Note that this action must be done as root, or it will have no effect; so a root-owned startup script is the easiest place to do it. If you do this, you should also set these environment variables in the startup script before invoking the postmaster:

```
export PG_OOM_ADJUST_FILE=/proc/self/oom_score_adj
export PG_OOM_ADJUST_VALUE=0
```

These settings will cause postmaster child processes to run with the normal OOM score adjustment of zero, so that the OOM killer can still target them at need. You could use some other value for `PG_OOM_ADJUST_VALUE` if you want the child processes to run with some other OOM score adjustment. (`PG_OOM_ADJUST_VALUE` can also be omitted, in which case it defaults to zero.) If you do not set `PG_OOM_ADJUST_FILE`, the child processes will run with the same OOM score adjustment as the postmaster, which is unwise since the whole point is to ensure that the postmaster has a preferential setting.

Older Linux kernels do not offer `/proc/self/oom_score_adj`, but may have a previous version of the same functionality called `/proc/self/oom_adj`. This works the same except the disable value is `-17` not `-1000`.

**Note:** Some vendors’ Linux 2.4 kernels are reported to have early versions of the 2.6 overcommit `sysctl` parameter. However, setting `vm.overcommit_memory` to 2 on a 2.4 kernel that does not have the relevant code will make things worse, not better. It is recommended that you inspect the actual kernel source code (see the function `vm_enough_memory` in the file `mm/mmap.c`) to verify what is supported in your kernel before you try this in a 2.4 installation. The presence of the `overcommit-accounting` documentation file should not be taken as evidence that the feature is there. If in any doubt, consult a kernel expert or your kernel vendor.

### 18.4.5. Linux Huge Pages

Using huge pages reduces overhead when using large contiguous chunks of memory, as PostgreSQL does, particularly when using large values of `shared_buffers`. To use this feature in PostgreSQL you need a kernel with `CONFIG_HUGETLBFS=y` and `CONFIG_HUGETLB_PAGE=y`. You will also have to adjust the kernel setting `vm.nr_hugepages`. To estimate the number of huge pages needed, start

\(^1\) http://lwn.net/Articles/104179/
Chapter 18. Server Setup and Operation

PostgreSQL without huge pages enabled and check the postmaster’s VmPeak value, as well as the system’s huge page size, using the /proc file system. This might look like:

```
$ head -1 $PGDATA/postmaster.pid
4170
$ grep ^VmPeak /proc/4170/status
VmPeak: 6490428 kB
$ grep ^Hugepagesize /proc/meminfo
Hugepagesize: 2048 kB
```

6490428 / 2048 gives approximately 3169.154, so in this example we need at least 3170 huge pages, which we can set with:

```
$ sysctl -w vm.nr_hugepages=3170
```

A larger setting would be appropriate if other programs on the machine also need huge pages. Don’t forget to add this setting to /etc/sysctl.conf so that it will be reapplied after reboots.

Sometimes the kernel is not able to allocate the desired number of huge pages immediately, so it might be necessary to repeat the command or to reboot. (Immediately after a reboot, most of the machine’s memory should be available to convert into huge pages.) To verify the huge page allocation situation, use:

```
$ grep Huge /proc/meminfo
```

It may also be necessary to give the database server’s operating system user permission to use huge pages by setting vm.hugetlb_shm_group via sysctl, and/or give permission to lock memory with ulimit -l.

The default behavior for huge pages in PostgreSQL is to use them when possible and to fall back to normal pages when failing. To enforce the use of huge pages, you can set huge_pages to on in postgresql.conf. Note that with this setting PostgreSQL will fail to start if not enough huge pages are available.


### 18.5. Shutting Down the Server

There are several ways to shut down the database server. You control the type of shutdown by sending different signals to the master postgres process.

**SIGTERM**

This is the Smart Shutdown mode. After receiving SIGTERM, the server disallows new connections, but lets existing sessions end their work normally. It shuts down only after all of the sessions terminate. If the server is in online backup mode, it additionally waits until online backup mode is no longer active. While backup mode is active, new connections will still be allowed, but only to superusers (this exception allows a superuser to connect to terminate online backup mode). If the server is in recovery when a smart shutdown is requested, recovery and streaming replication will be stopped only after all regular sessions have terminated.
Chapter 18. Server Setup and Operation

SIGINT

This is the Fast Shutdown mode. The server disallows new connections and sends all existing server processes SIGTERM, which will cause them to abort their current transactions and exit promptly. It then waits for all server processes to exit and finally shuts down. If the server is in online backup mode, backup mode will be terminated, rendering the backup useless.

SIGQUIT

This is the Immediate Shutdown mode. The server will send SIGQUIT to all child processes and wait for them to terminate. If any do not terminate within 5 seconds, they will be sent SIGKILL. The master server process exits as soon as all child processes have exited, without doing normal database shutdown processing. This will lead to recovery (by replaying the WAL log) upon next start-up. This is recommended only in emergencies.

The pg_ctl program provides a convenient interface for sending these signals to shut down the server. Alternatively, you can send the signal directly using `kill` on non-Windows systems. The PID of the postgres process can be found using the `ps` program, or from the file `postmaster.pid` in the data directory. For example, to do a fast shutdown:

```
$ kill -INT `head -1 /usr/local/pgsql/data/postmaster.pid`
```

**Important:** It is best not to use SIGKILL to shut down the server. Doing so will prevent the server from releasing shared memory and semaphores, which might then have to be done manually before a new server can be started. Furthermore, SIGKILL kills the postgres process without letting it relay the signal to its subprocesses, so it will be necessary to kill the individual subprocesses by hand as well.

To terminate an individual session while allowing other sessions to continue, use `pg_terminate_backend()` (see Table 9-77) or send a SIGTERM signal to the child process associated with the session.

18.6. Upgrading a PostgreSQL Cluster

This section discusses how to upgrade your database data from one PostgreSQL release to a newer one.

PostgreSQL major versions are represented by the first two digit groups of the version number, e.g., 8.4. PostgreSQL minor versions are represented by the third group of version digits, e.g., 8.4.2 is the second minor release of 8.4. Minor releases never change the internal storage format and are always compatible with earlier and later minor releases of the same major version number, e.g., 8.4.2 is compatible with 8.4, 8.4.1 and 8.4.6. To update between compatible versions, you simply replace the executables while the server is down and restart the server. The data directory remains unchanged — minor upgrades are that simple.

For major releases of PostgreSQL, the internal data storage format is subject to change, thus complicating upgrades. The traditional method for moving data to a new major version is to dump and reload the database, though this can be slow. A faster method is `pg_upgrade`. Replication methods are also available, as discussed below.
New major versions also typically introduce some user-visible incompatibilities, so application programming changes might be required. All user-visible changes are listed in the release notes (Appendix E); pay particular attention to the section labeled "Migration". If you are upgrading across several major versions, be sure to read the release notes for each intervening version.

Cautious users will want to test their client applications on the new version before switching over fully; therefore, it’s often a good idea to set up concurrent installations of old and new versions. When testing a PostgreSQL major upgrade, consider the following categories of possible changes:

Administration

The capabilities available for administrators to monitor and control the server often change and improve in each major release.

SQL

Typically this includes new SQL command capabilities and not changes in behavior, unless specifically mentioned in the release notes.

Library API

Typically libraries like libpq only add new functionality, again unless mentioned in the release notes.

System Catalogs

System catalog changes usually only affect database management tools.

Server C-language API

This involves changes in the backend function API, which is written in the C programming language. Such changes affect code that references backend functions deep inside the server.

18.6.1. Upgrading Data via pg_dumpall

One upgrade method is to dump data from one major version of PostgreSQL and reload it in another — to do this, you must use a logical backup tool like pg_dumpall; file system level backup methods will not work. (There are checks in place that prevent you from using a data directory with an incompatible version of PostgreSQL, so no great harm can be done by trying to start the wrong server version on a data directory.)

It is recommended that you use the pg_dump and pg_dumpall programs from the newer version of PostgreSQL, to take advantage of enhancements that might have been made in these programs. Current releases of the dump programs can read data from any server version back to 7.0.

These instructions assume that your existing installation is under the /usr/local/pgsql directory, and that the data area is in /usr/local/pgsql/data. Substitute your paths appropriately.

1. If making a backup, make sure that your database is not being updated. This does not affect the integrity of the backup, but the changed data would of course not be included. If necessary, edit the permissions in the file /usr/local/pgsql/data/pg_hba.conf (or equivalent) to disallow access from everyone except you. See Chapter 20 for additional information on access control.

   To back up your database installation, type:
   
   pg_dumpall > outputfile
   
   To make the backup, you can use the pg_dumpall command from the version you are currently running; see Section 25.1.2 for more details. For best results, however, try to use the pg_dumpall
command from PostgreSQL 9.6.17, since this version contains bug fixes and improvements over older versions. While this advice might seem idiosyncratic since you haven’t installed the new version yet, it is advisable to follow it if you plan to install the new version in parallel with the old version. In that case you can complete the installation normally and transfer the data later. This will also decrease the downtime.

2. Shut down the old server:
   ```bash
   pg_ctl stop
   ```
   On systems that have PostgreSQL started at boot time, there is probably a start-up file that will accomplish the same thing. For example, on a Red Hat Linux system one might find that this works:
   ```bash
   /etc/rc.d/init.d/postgresql stop
   ```
   See Chapter 18 for details about starting and stopping the server.

3. If restoring from backup, rename or delete the old installation directory if it is not version-specific. It is a good idea to rename the directory, rather than delete it, in case you have trouble and need to revert to it. Keep in mind the directory might consume significant disk space. To rename the directory, use a command like this:
   ```bash
   mv /usr/local/pgsql /usr/local/pgsql.old
   ```
   (Be sure to move the directory as a single unit so relative paths remain unchanged.)

4. Install the new version of PostgreSQL as outlined in Section 16.4.

5. Create a new database cluster if needed. Remember that you must execute these commands while logged in to the special database user account (which you already have if you are upgrading).
   ```bash
   /usr/local/pgsql/bin/initdb -D /usr/local/pgsql/data
   ```

6. Restore your previous `pg_hba.conf` and any `postgresql.conf` modifications.

7. Start the database server, again using the special database user account:
   ```bash
   /usr/local/pgsql/bin/postgres -D /usr/local/pgsql/data
   ```

8. Finally, restore your data from backup with:
   ```bash
   /usr/local/pgsql/bin/psql -d postgres -f outputfile
   ```
   using the new `psql`.

The least downtime can be achieved by installing the new server in a different directory and running both the old and the new servers in parallel, on different ports. Then you can use something like:

```bash
pg_dumpall -p 5432 | psql -d postgres -p 5433
```

to transfer your data.

### 18.6.2. Upgrading Data via pg_upgrade

The `pg_upgrade` module allows an installation to be migrated in-place from one major PostgreSQL version to another. Upgrades can be performed in minutes, particularly with `--link` mode. It requires steps similar to `pg_dumpall` above, e.g. starting/stopping the server, running `initdb`. The `pg_upgrade` documentation outlines the necessary steps.
18.6.3. Upgrading Data via Replication

It is also possible to use certain replication methods, such as Slony, to create a standby server with the updated version of PostgreSQL. This is possible because Slony supports replication between different major versions of PostgreSQL. The standby can be on the same computer or a different computer. Once it has synced up with the master server (running the older version of PostgreSQL), you can switch masters and make the standby the master and shut down the older database instance. Such a switch-over results in only several seconds of downtime for an upgrade.

18.7. Preventing Server Spoofing

While the server is running, it is not possible for a malicious user to take the place of the normal database server. However, when the server is down, it is possible for a local user to spoof the normal server by starting their own server. The spoof server could read passwords and queries sent by clients, but could not return any data because the $PGDATA$ directory would still be secure because of directory permissions. Spoofing is possible because any user can start a database server; a client cannot identify an invalid server unless it is specially configured.

One way to prevent spoofing of local connections is to use a Unix domain socket directory (unix_socket_directories) that has write permission only for a trusted local user. This prevents a malicious user from creating their own socket file in that directory. If you are concerned that some applications might still reference /tmp for the socket file and hence be vulnerable to spoofing, during operating system startup create a symbolic link /tmp/.s.PGSQL.5432 that points to the relocated socket file. You also might need to modify your /tmp cleanup script to prevent removal of the symbolic link.

Another option for local connections is for clients to use requirepeer to specify the required owner of the server process connected to the socket.

To prevent spoofing on TCP connections, the best solution is to use SSL certificates and make sure that clients check the server’s certificate. To do that, the server must be configured to accept only hostssl connections (Section 20.1) and have SSL key and certificate files (Section 18.9). The TCP client must connect using sslmode=verify-ca or verify-full and have the appropriate root certificate file installed (Section 32.18.1).

18.8. Encryption Options

PostgreSQL offers encryption at several levels, and provides flexibility in protecting data from disclosure due to database server theft, unscrupulous administrators, and insecure networks. Encryption might also be required to secure sensitive data such as medical records or financial transactions.

Password Storage Encryption

By default, database user passwords are stored as MD5 hashes, so the administrator cannot determine the actual password assigned to the user. If MD5 encryption is used for client authentication, the unencrypted password is never even temporarily present on the server because the client MD5-encrypts it before being sent across the network.

Encryption For Specific Columns

The pgcrypto module allows certain fields to be stored encrypted. This is useful if only some of the data is sensitive. The client supplies the decryption key and the data is decrypted on the
server and then sent to the client.

The decrypted data and the decryption key are present on the server for a brief time while it is being decrypted and communicated between the client and server. This presents a brief moment where the data and keys can be intercepted by someone with complete access to the database server, such as the system administrator.

Data Partition Encryption

Storage encryption can be performed at the file system level or the block level. Linux file system encryption options include eCryptfs and EncFS, while FreeBSD uses PEFS. Block level or full disk encryption options include dm-crypt + LUKS on Linux and GEOM modules geli and gbde on FreeBSD. Many other operating systems support this functionality, including Windows.

This mechanism prevents unencrypted data from being read from the drives if the drives or the entire computer is stolen. This does not protect against attacks while the file system is mounted, because when mounted, the operating system provides an unencrypted view of the data. However, to mount the file system, you need some way for the encryption key to be passed to the operating system, and sometimes the key is stored somewhere on the host that mounts the disk.

Encrypting Passwords Across A Network

The MD5 authentication method double-encrypts the password on the client before sending it to the server. It first MD5-encrypts it based on the user name, and then encrypts it based on a random salt sent by the server when the database connection was made. It is this double-encrypted value that is sent over the network to the server. Double-encryption not only prevents the password from being discovered, it also prevents another connection from using the same encrypted password to connect to the database server at a later time.

Encrypting Data Across A Network

SSL connections encrypt all data sent across the network: the password, the queries, and the data returned. The pg_hba.conf file allows administrators to specify which hosts can use non-encrypted connections (host) and which require SSL-encrypted connections (hostssl). Also, clients can specify that they connect to servers only via SSL. Stunnel or SSH can also be used to encrypt transmissions.

SSL Host Authentication

It is possible for both the client and server to provide SSL certificates to each other. It takes some extra configuration on each side, but this provides stronger verification of identity than the mere use of passwords. It prevents a computer from pretending to be the server just long enough to read the password sent by the client. It also helps prevent “man in the middle” attacks where a computer between the client and server pretends to be the server and reads and passes all data between the client and server.

Client-Side Encryption

If the system administrator for the server’s machine cannot be trusted, it is necessary for the client to encrypt the data; this way, unencrypted data never appears on the database server. Data is encrypted on the client before being sent to the server, and database results have to be decrypted on the client before being used.

18.9. Secure TCP/IP Connections with SSL

PostgreSQL has native support for using SSL connections to encrypt client/server communications for increased security. This requires that OpenSSL is installed on both client and server systems and
Chapter 18. Server Setup and Operation

that support in PostgreSQL is enabled at build time (see Chapter 16).

With SSL support compiled in, the PostgreSQL server can be started with SSL enabled by setting the
parameter ssl to on in postgresql.conf. The server will listen for both normal and SSL connections
on the same TCP port, and will negotiate with any connecting client on whether to use SSL. By
default, this is at the client’s option; see Section 20.1 about how to set up the server to require use of
SSL for some or all connections.

PostgreSQL reads the system-wide OpenSSL configuration file. By default, this file is named
openssl.cnf and is located in the directory reported by openssl version -d. This default
can be overridden by setting environment variable OPENSSL_CONF to the name of the desired
configuration file.

OpenSSL supports a wide range of ciphers and authentication algorithms, of varying strength. While a
list of ciphers can be specified in the OpenSSL configuration file, you can specify ciphers specifically
for use by the database server by modifying ssl_ciphers in postgresql.conf.

Note: It is possible to have authentication without encryption overhead by using NULL-SHA or
NULL-MD5 ciphers. However, a man-in-the-middle could read and pass communications between
client and server. Also, encryption overhead is minimal compared to the overhead of authentica-
tion. For these reasons NULL ciphers are not recommended.

To start in SSL mode, files containing the server certificate and private key must exist. By default,
these files are expected to be named server.crt and server.key, respectively, in the server’s
data directory, but other names and locations can be specified using the configuration parameters
ssl_cert_file and ssl_key_file.

On Unix systems, the permissions on server.key must disallow any access to world or group;
achieve this by the command chmod 0600 server.key. Alternatively, the file can be owned by root
and have group read access (that is, 0640 permissions). That setup is intended for installations where
certificate and key files are managed by the operating system. The user under which the PostgreSQL
server runs should then be made a member of the group that has access to those certificate and key
files.

If the private key is protected with a passphrase, the server will prompt for the passphrase and will
not start until it has been entered.

The first certificate in server.crt must be the server’s certificate because it must match the server’s
private key. The certificates of “intermediate” certificate authorities can also be appended to the file.
Doing this avoids the necessity of storing intermediate certificates on clients, assuming the root and
intermediate certificates were created with v3_ca extensions. This allows easier expiration of inter-
mediate certificates.

It is not necessary to add the root certificate to server.crt. Instead, clients must have the root
certificate of the server’s certificate chain.

18.9.1. Using Client Certificates

To require the client to supply a trusted certificate, place certificates of the root certificate authorities
(CAs) you trust in a file in the data directory, set the parameter ssl_ca_file in postgresql.conf
to the new file name, and add the authentication option clientcert=1 to the appropriate hostssl
line(s) in pg_hba.conf. A certificate will then be requested from the client during SSL connection
startup. (See Section 32.18 for a description of how to set up certificates on the client.) The server
will verify that the client’s certificate is signed by one of the trusted certificate authorities.
Intermediate certificates that chain up to existing root certificates can also appear in the file `root.crt` if you wish to avoid storing them on clients (assuming the root and intermediate certificates were created with `v3_ca` extensions). Certificate Revocation List (CRL) entries are also checked if the parameter `ssl_crl_file` is set. (See http://h41379.www4.hpe.com/doc/83final/ba554_90007/ch04s02.html for diagrams showing SSL certificate usage.)

The `clientcert` authentication option is available for all authentication methods, but only in `pg_hba.conf` lines specified as `hostssl`. When `clientcert` is not specified or is set to 0, the server will still verify any presented client certificates against its CA file, if one is configured — but it will not insist that a client certificate be presented.

If you are setting up client certificates, you may wish to use the `cert` authentication method, so that the certificates control user authentication as well as providing connection security. See Section 20.3.9 for details. (It is not necessary to specify `clientcert=1` explicitly when using the `cert` authentication method.)

### 18.9.2. SSL Server File Usage

Table 18-2 summarizes the files that are relevant to the SSL setup on the server. (The shown file names are default or typical names. The locally configured names could be different.)

<table>
<thead>
<tr>
<th>File</th>
<th>Contents</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ssl_cert_file</code></td>
<td>server certificate</td>
<td>sent to client to indicate server’s identity</td>
</tr>
<tr>
<td>(<code>$PGDATA/server.crt</code>)</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ssl_key_file</code></td>
<td>server private key</td>
<td>proves server certificate was sent by the owner; does not indicate certificate owner is trustworthy</td>
</tr>
<tr>
<td>(<code>$PGDATA/server.key</code>)</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ssl_ca_file</code></td>
<td>trusted certificate authorities</td>
<td>checks that client certificate is signed by a trusted certificate authority</td>
</tr>
<tr>
<td>(<code>$PGDATA/root.crt</code>)</td>
<td></td>
<td></td>
</tr>
<tr>
<td><code>ssl_crl_file</code></td>
<td>certificates revoked by certificate authorities</td>
<td>client certificate must not be on this list</td>
</tr>
<tr>
<td>(<code>$PGDATA/root.crl</code>)</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

The files `server.key`, `server.crt`, `root.crt`, and `root.crl` (or their configured alternative names) are only examined during server start; so you must restart the server for changes in them to take effect.

### 18.9.3. Creating Certificates

To create a simple self-signed certificate for the server, valid for 365 days, use the following OpenSSL command, replacing `dbhost.yourdomain.com` with the server’s host name:

```
openssl req -new -x509 -days 365 -nodes -text -out server.crt \
   -keyout server.key -subj "/CN=dbhost.yourdomain.com"
```

Then do:

```
chmod og-rwx server.key
```
Chapter 18. Server Setup and Operation

because the server will reject the file if its permissions are more liberal than this. For more details on how to create your server private key and certificate, refer to the OpenSSL documentation.

While a self-signed certificate can be used for testing, a certificate signed by a certificate authority (CA) (usually an enterprise-wide root CA) should be used in production.

To create a server certificate whose identity can be validated by clients, first create a certificate signing request (CSR) and a public/private key file:

```
openssl req -new -nodes -text -out root.csr \
    -keyout root.key -subj "/CN=root.yourdomain.com"
chmod og-rwx root.key
```

Then, sign the request with the key to create a root certificate authority (using the default OpenSSL configuration file location on Linux):

```
openssl x509 -req -in root.csr -text -days 3650 \
    -extfile /etc/ssl/openssl.cnf -extensions v3_ca \
    -signkey root.key -out root.crt
```

Finally, create a server certificate signed by the new root certificate authority:

```
openssl req -new -nodes -text -out server.csr \
    -keyout server.key -subj "/CN=dbhost.yourdomain.com"
chmod og-rwx server.key
openssl x509 -req -in server.csr -text -days 365 \
    -CA root.crt -CAkey root.key -CAcreateserial \
    -out server.crt
```

server.crt and server.key should be stored on the server, and root.crt should be stored on the client so the client can verify that the server’s leaf certificate was signed by its trusted root certificate. root.key should be stored offline for use in creating future certificates.

It is also possible to create a chain of trust that includes intermediate certificates:

```
# root
openssl req -new -nodes -text -out root.csr \
    -keyout root.key -subj "/CN=root.yourdomain.com"
chmod og-rwx root.key
openssl x509 -req -in root.csr -text -days 3650 \
    -extfile /etc/ssl/openssl.cnf -extensions v3_ca \
    -signkey root.key -out root.crt
```

```
# intermediate
openssl req -new -nodes -text -out intermediate.csr \
    -keyout intermediate.key -subj "/CN=intermediate.yourdomain.com"
chmod og-rwx intermediate.key
openssl x509 -req -in intermediate.csr -text -days 1825 \
    -extfile /etc/ssl/openssl.cnf -extensions v3_ca \
    -CA root.crt -CAkey root.key -CAcreateserial \
    -out intermediate.crt
```

```
# leaf
openssl req -new -nodes -text -out server.csr \
    -keyout server.key -subj "/CN=dbhost.yourdomain.com"
chmod og-rwx server.key
openssl x509 -req -in server.csr -text -days 365 \
    -CA intermediate.crt -CAkey intermediate.key -CAcreateserial \
```
server.crt and intermediate.crt should be concatenated into a certificate file bundle and stored on the server. server.key should also be stored on the server. root.crt should be stored on the client so the client can verify that the server’s leaf certificate was signed by a chain of certificates linked to its trusted root certificate. root.key and intermediate.key should be stored offline for use in creating future certificates.

18.10. Secure TCP/IP Connections with SSH Tunnels

It is possible to use SSH to encrypt the network connection between clients and a PostgreSQL server. Done properly, this provides an adequately secure network connection, even for non-SSL-capable clients.

First make sure that an SSH server is running properly on the same machine as the PostgreSQL server and that you can log in using ssh as some user. Then you can establish a secure tunnel with a command like this from the client machine:

```
ssh -L 63333:localhost:5432 joe@foo.com
```

The first number in the -L argument, 63333, is the port number of your end of the tunnel; it can be any unused port. (IANA reserves ports 49152 through 65535 for private use.) The second number, 5432, is the remote end of the tunnel: the port number your server is using. The name or IP address between the port numbers is the host with the database server you are going to connect to, as seen from the host you are logging in to, which is foo.com in this example. In order to connect to the database server using this tunnel, you connect to port 63333 on the local machine:

```
psql -h localhost -p 63333 postgres
```

To the database server it will then look as though you are really user joe on host foo.com connecting to localhost in that context, and it will use whatever authentication procedure was configured for connections from this user and host. Note that the server will not think the connection is SSL-encrypted, since in fact it is not encrypted between the SSH server and the PostgreSQL server. This should not pose any extra security risk as long as they are on the same machine.

In order for the tunnel setup to succeed you must be allowed to connect via ssh as joe@foo.com, just as if you had attempted to use ssh to create a terminal session.

You could also have set up the port forwarding as

```
ssh -L 63333:foo.com:5432 joe@foo.com
```

but then the database server will see the connection as coming in on its foo.com interface, which is not opened by the default setting listen_addresses = ‘localhost’. This is usually not what you want.

If you have to “hop” to the database server via some login host, one possible setup could look like this:

```
ssh -L 63333:db.foo.com:5432 joe@shell.foo.com
```

Note that this way the connection from shell.foo.com to db.foo.com will not be encrypted by the SSH tunnel. SSH offers quite a few configuration possibilities when the network is restricted in various ways. Please refer to the SSH documentation for details.
Tip: Several other applications exist that can provide secure tunnels using a procedure similar in concept to the one just described.

18.11. Registering Event Log on Windows

To register a Windows event log library with the operating system, issue this command:

```
regsvr32 pgsql_library_directory/pgevent.dll
```

This creates registry entries used by the event viewer, under the default event source named PostgreSQL.

To specify a different event source name (see event_source), use the /n and /i options:

```
regsvr32 /n /i:event_source_name pgsql_library_directory/pgevent.dll
```

To unregister the event log library from the operating system, issue this command:

```
regsvr32 /u [/i:event_source_name] pgsql_library_directory/pgevent.dll
```

Note: To enable event logging in the database server, modify log_destination to include eventlog in postgresql.conf.
Chapter 19. Server Configuration

There are many configuration parameters that affect the behavior of the database system. In the first section of this chapter we describe how to interact with configuration parameters. The subsequent sections discuss each parameter in detail.

19.1. Setting Parameters

19.1.1. Parameter Names and Values

All parameter names are case-insensitive. Every parameter takes a value of one of five types: boolean, string, integer, floating point, or enumerated (enum). The type determines the syntax for setting the parameter:

- **Boolean**: Values can be written as on, off, true, false, yes, no, 1, 0 (all case-insensitive) or any unambiguous prefix of one of these.
- **String**: In general, enclose the value in single quotes, doubling any single quotes within the value. Quotes can usually be omitted if the value is a simple number or identifier, however.
- **Numeric (integer and floating point)**: A decimal point is permitted only for floating-point parameters. Do not use thousands separators. Quotes are not required.
- **Numeric with Unit**: Some numeric parameters have an implicit unit, because they describe quantities of memory or time. The unit might be kilobytes, blocks (typically eight kilobytes), milliseconds, seconds, or minutes. An unadorned numeric value for one of these settings will use the setting’s default unit, which can be learned from pg_settings.unit. For convenience, settings can be given with a unit specified explicitly, for example ‘120 ms’ for a time value, and they will be converted to whatever the parameter’s actual unit is. Note that the value must be written as a string (with quotes) to use this feature. The unit name is case-sensitive, and there can be whitespace between the numeric value and the unit.
  - Valid memory units are kB (kilobytes), MB (megabytes), GB (gigabytes), and TB (terabytes). The multiplier for memory units is 1024, not 1000.
  - Valid time units are ms (milliseconds), s (seconds), min (minutes), h (hours), and d (days).
- **Enumerated**: Enumerated-type parameters are written in the same way as string parameters, but are restricted to have one of a limited set of values. The values allowable for such a parameter can be found from pg_settings enumvals. Enum parameter values are case-insensitive.

19.1.2. Parameter Interaction via the Configuration File

The most fundamental way to set these parameters is to edit the file postgresql.conf, which is normally kept in the data directory. A default copy is installed when the database cluster directory is initialized. An example of what this file might look like is:

```plaintext
# This is a comment
log_connections = yes
log_destination = 'syslog'
search_path = "'$user', public'
shared_buffers = 128MB
```
One parameter is specified per line. The equal sign between name and value is optional. Whitespace is insignificant (except within a quoted parameter value) and blank lines are ignored. Hash marks (#) designate the remainder of the line as a comment. Parameter values that are not simple identifiers or numbers must be single-quoted. To embed a single quote in a parameter value, write either two quotes (preferred) or backslash-quote. If the file contains multiple entries for the same parameter, all but the last one are ignored.

Parameters set in this way provide default values for the cluster. The settings seen by active sessions will be these values unless they are overridden. The following sections describe ways in which the administrator or user can override these defaults.

The configuration file is reread whenever the main server process receives a SIGHUP signal; this signal is most easily sent by running `pg_ctl reload` from the command line or by calling the SQL function `pg_reload_conf()`. The main server process also propagates this signal to all currently running server processes, so that existing sessions also adopt the new values (this will happen after they complete any currently-executing client command). Alternatively, you can send the signal to a single server process directly. Some parameters can only be set at server start; any changes to their entries in the configuration file will be ignored until the server is restarted. Invalid parameter settings in the configuration file are likewise ignored (but logged) during SIGHUP processing.

In addition to `postgresql.conf`, a PostgreSQL data directory contains a file `postgresql.auto.conf`, which has the same format as `postgresql.conf` but is intended to be edited automatically not manually. This file holds settings provided through the ALTER SYSTEM command. This file is read whenever `postgresql.conf` is, and its settings take effect in the same way. Settings in `postgresql.auto.conf` override those in `postgresql.conf`.

External tools may also modify `postgresql.auto.conf`. It is not recommended to do this while the server is running, since a concurrent ALTER SYSTEM command could overwrite such changes. Such tools might simply append new settings to the end, or they might choose to remove duplicate settings and/or comments (as ALTER SYSTEM will).

The system view `pg_file_settings` can be helpful for pre-testing changes to the configuration files, or for diagnosing problems if a SIGHUP signal did not have the desired effects.

### 19.1.3. Parameter Interaction via SQL

PostgreSQL provides three SQL commands to establish configuration defaults. The already-mentioned ALTER SYSTEM command provides a SQL-accessible means of changing global defaults; it is functionally equivalent to editing `postgresql.conf`. In addition, there are two commands that allow setting of defaults on a per-database or per-role basis:

- The ALTER DATABASE command allows global settings to be overridden on a per-database basis.
- The ALTER ROLE command allows both global and per-database settings to be overridden with user-specific values.

Values set with ALTER DATABASE and ALTER ROLE are applied only when starting a fresh database session. They override values obtained from the configuration files or server command line, and constitute defaults for the rest of the session. Note that some settings cannot be changed after server start, and so cannot be set with these commands (or the ones listed below).

Once a client is connected to the database, PostgreSQL provides two additional SQL commands (and equivalent functions) to interact with session-local configuration settings;
Chapter 19. Server Configuration

• The SHOW command allows inspection of the current value of all parameters. The corresponding function is `current_setting(setting_name text)`.

• The SET command allows modification of the current value of those parameters that can be set locally to a session; it has no effect on other sessions. The corresponding function is `set_config(setting_name, new_value, is_local)`.

In addition, the system view `pg_settings` can be used to view and change session-local values:

• Querying this view is similar to using `SHOW ALL` but provides more detail. It is also more flexible, since it’s possible to specify filter conditions or join against other relations.

• Using UPDATE on this view, specifically updating the `setting` column, is the equivalent of issuing SET commands. For example, the equivalent of

```
SET configuration_parameter TO DEFAULT;
```

is:

```
UPDATE pg_settings SET setting = reset_val WHERE name = 'configuration_parameter';
```

19.1.4. Parameter Interaction via the Shell

In addition to setting global defaults or attaching overrides at the database or role level, you can pass settings to PostgreSQL via shell facilities. Both the server and libpq client library accept parameter values via the shell.

• During server startup, parameter settings can be passed to the `postgres` command via the `-c` command-line parameter. For example,

```
postgres -c log_connections=yes -c log_destination=syslog
```

Settings provided in this way override those set via `postgresql.conf` or `ALTER SYSTEM`, so they cannot be changed globally without restarting the server.

• When starting a client session via libpq, parameter settings can be specified using the `PGOPTIONS` environment variable. Settings established in this way constitute defaults for the life of the session, but do not affect other sessions. For historical reasons, the format of `PGOPTIONS` is similar to that used when launching the `postgres` command; specifically, the `-c` flag must be specified. For example,

```
env PGOPTIONS="-c geqo=off -c statement_timeout=5min" psql
```

Other clients and libraries might provide their own mechanisms, via the shell or otherwise, that allow the user to alter session settings without direct use of SQL commands.

19.1.5. Managing Configuration File Contents

PostgreSQL provides several features for breaking down complex `postgresql.conf` files into sub-files. These features are especially useful when managing multiple servers with related, but not identical, configurations.

In addition to individual parameter settings, the `postgresql.conf` file can contain `include` directives, which specify another file to read and process as if it were inserted into the configuration file at this point. This feature allows a configuration file to be divided into physically separate parts. Include directives simply look like:

```
include ‘filename’
```
If the file name is not an absolute path, it is taken as relative to the directory containing the referencing configuration file. Inclusions can be nested.

There is also an `include_if_exists` directive, which acts the same as the `include` directive, except when the referenced file does not exist or cannot be read. A regular `include` will consider this an error condition, but `include_if_exists` merely logs a message and continues processing the referencing configuration file.

The `postgresql.conf` file can also contain `include_dir` directives, which specify an entire directory of configuration files to include. These look like

```plaintext
include_dir 'directory'
```

Non-absolute directory names are taken as relative to the directory containing the referencing configuration file. Within the specified directory, only non-directory files whose names end with the suffix `.conf` will be included. File names that start with the `.` character are also ignored, to prevent mistakes since such files are hidden on some platforms. Multiple files within an include directory are processed in file name order (according to C locale rules, i.e. numbers before letters, and uppercase letters before lowercase ones).

Include files or directories can be used to logically separate portions of the database configuration, rather than having a single large `postgresql.conf` file. Consider a company that has two database servers, each with a different amount of memory. There are likely elements of the configuration both will share, for things such as logging. But memory-related parameters on the server will vary between the two. And there might be server specific customizations, too. One way to manage this situation is to break the custom configuration changes for your site into three files. You could add this to the end of your `postgresql.conf` file to include them:

```plaintext
include 'shared.conf'
include 'memory.conf'
include 'server.conf'
```

All systems would have the same `shared.conf`. Each server with a particular amount of memory could share the same `memory.conf`; you might have one for all servers with 8GB of RAM, another for those having 16GB. And finally `server.conf` could have truly server-specific configuration information in it.

Another possibility is to create a configuration file directory and put this information into files there. For example, a `conf.d` directory could be referenced at the end of `postgresql.conf`:

```plaintext
include_dir 'conf.d'
```

Then you could name the files in the `conf.d` directory like this:

```plaintext
00shared.conf
01memory.conf
02server.conf
```

This naming convention establishes a clear order in which these files will be loaded. This is important because only the last setting encountered for a particular parameter while the server is reading configuration files will be used. In this example, something set in `conf.d/02server.conf` would override a value set in `conf.d/01memory.conf`.

You might instead use this approach to naming the files descriptively:

```plaintext
00shared.conf
01memory-8GB.conf
02server-foo.conf
```
Chapter 19. Server Configuration

This sort of arrangement gives a unique name for each configuration file variation. This can help eliminate ambiguity when several servers have their configurations all stored in one place, such as in a version control repository. (Storing database configuration files under version control is another good practice to consider.)

19.2. File Locations

In addition to the `postgresql.conf` file already mentioned, PostgreSQL uses two other manually-edited configuration files, which control client authentication (their use is discussed in Chapter 20). By default, all three configuration files are stored in the database cluster's data directory. The parameters described in this section allow the configuration files to be placed elsewhere. (Doing so can ease administration. In particular it is often easier to ensure that the configuration files are properly backed-up when they are kept separate.)

data_directory(string)

Specifies the directory to use for data storage. This parameter can only be set at server start.

config_file(string)

Specifies the main server configuration file (customarily called `postgresql.conf`). This parameter can only be set on the `postgres` command line.

hba_file(string)

Specifies the configuration file for host-based authentication (customarily called `pg_hba.conf`). This parameter can only be set at server start.

ident_file(string)

Specifies the configuration file for user name mapping (customarily called `pg_ident.conf`). This parameter can only be set at server start. See also Section 20.2.

external_pid_file(string)

Specifies the name of an additional process-ID (PID) file that the server should create for use by server administration programs. This parameter can only be set at server start.

In a default installation, none of the above parameters are set explicitly. Instead, the data directory is specified by the `-D` command-line option or the `PGDATA` environment variable, and the configuration files are all found within the data directory.

If you wish to keep the configuration files elsewhere than the data directory, the `postgres -D` command-line option or `PGDATA` environment variable must point to the directory containing the configuration files, and the `data_directory` parameter must be set in `postgresql.conf` (or on the command line) to show where the data directory is actually located. Notice that `data_directory` overrides `-D` and `PGDATA` for the location of the data directory, but not for the location of the configuration files.

If you wish, you can specify the configuration file names and locations individually using the parameters `config_file`, `hba_file` and/or `ident_file`. `config_file` can only be specified on the `postgres` command line, but the others can be set within the main configuration file. If all three parameters plus `data_directory` are explicitly set, then it is not necessary to specify `-D` or `PGDATA`.

When setting any of these parameters, a relative path will be interpreted with respect to the directory in which `postgres` is started.
19.3. Connections and Authentication

19.3.1. Connection Settings

**listen_addresses**: (string)

Specifies the TCP/IP address(es) on which the server is to listen for connections from client applications. The value takes the form of a comma-separated list of host names and/or numeric IP addresses. The special entry * corresponds to all available IP interfaces. The entry 0.0.0.0 allows listening for all IPv4 addresses and :: allows listening for all IPv6 addresses. If the list is empty, the server does not listen on any IP interface at all, in which case only Unix-domain sockets can be used to connect to it. The default value is localhost, which allows only local TCP/IP “loopback” connections to be made. While client authentication (Chapter 20) allows fine-grained control over who can access the server, listen_addresses controls which interfaces accept connection attempts, which can help prevent repeated malicious connection requests on insecure network interfaces. This parameter can only be set at server start.

**port**: (integer)

The TCP port the server listens on; 5432 by default. Note that the same port number is used for all IP addresses the server listens on. This parameter can only be set at server start.

**max_connections**: (integer)

Determines the maximum number of concurrent connections to the database server. The default is typically 100 connections, but might be less if your kernel settings will not support it (as determined during initdb). This parameter can only be set at server start.

When running a standby server, you must set this parameter to the same or higher value than on the master server. Otherwise, queries will not be allowed in the standby server.

**superuser_reserved_connections**: (integer)

Determines the number of connection “slots” that are reserved for connections by PostgreSQL superusers. At most max_connections connections can ever be active simultaneously. Whenever the number of active concurrent connections is at least max_connections minus superuser_reserved_connections, new connections will be accepted only for superusers, and no new replication connections will be accepted.

The default value is three connections. The value must be less than the value of max_connections. This parameter can only be set at server start.

**unix_socket_directories**: (string)

Specifies the directory of the Unix-domain socket(s) on which the server is to listen for connections from client applications. Multiple sockets can be created by listing multiple directories separated by commas. Whitespace between entries is ignored; surround a directory name with double quotes if you need to include whitespace or commas in the name. An empty value specifies not listening on any Unix-domain sockets, in which case only TCP/IP sockets can be used to connect to the server. The default value is normally /tmp, but that can be changed at build time. This parameter can only be set at server start.

In addition to the socket file itself, which is named .s.PGSQL.nnnn where nnnn is the server’s port number, an ordinary file named .s.PGSQL.nnnn.lock will be created in each of the unix_socket_directories directories. Neither file should ever be removed manually.

This parameter is irrelevant on Windows, which does not have Unix-domain sockets.
unix_socket_group (string)

Sets the owning group of the Unix-domain socket(s). (The owning user of the sockets is always
the user that starts the server.) In combination with the parameter unix_socket_permissions
this can be used as an additional access control mechanism for Unix-domain connections. By
default this is the empty string, which uses the default group of the server user. This parameter
can only be set at server start.

This parameter is irrelevant on Windows, which does not have Unix-domain sockets.

unix_socket_permissions (integer)

Sets the access permissions of the Unix-domain socket(s). Unix-domain sockets use the usual
Unix file system permission set. The parameter value is expected to be a numeric mode specified
in the format accepted by the chmod and umask system calls. (To use the customary octal format
the number must start with a 0 (zero).)

The default permissions are 0777, meaning anyone can connect. Reasonable alternatives are
0770 (only user and group, see also unix_socket_group) and 0700 (only user). (Note that for
a Unix-domain socket, only write permission matters, so there is no point in setting or revoking
read or execute permissions.)

This access control mechanism is independent of the one described in Chapter 20.

This parameter can only be set at server start.

This parameter is irrelevant on systems, notably Solaris as of Solaris 10, that ignore
socket permissions entirely. There, one can achieve a similar effect by pointing
unix_socket_directories to a directory having search permission limited to the desired
audience. This parameter is also irrelevant on Windows, which does not have Unix-domain
sockets.

bonjour (boolean)

Enables advertising the server’s existence via Bonjour. The default is off. This parameter can
only be set at server start.

bonjour_name (string)

Specifies the Bonjour service name. The computer name is used if this parameter is set to the
empty string " (which is the default). This parameter is ignored if the server was not compiled
with Bonjour support. This parameter can only be set at server start.

tcp_keepalives_idle (integer)

Specifies the number of seconds of inactivity after which TCP should send a keepalive message
to the client. A value of 0 uses the system default. This parameter is supported only on systems
that support TCP_KEEPIDLE or an equivalent socket option, and on Windows; on other systems,
it must be zero. In sessions connected via a Unix-domain socket, this parameter is ignored and
always reads as zero.

Note: On Windows, a value of 0 will set this parameter to 2 hours, since Windows does not
provide a way to read the system default value.

tcp_keepalives_interval (integer)

Specifies the number of seconds after which a TCP keepalive message that is not acknowledged
by the client should be retransmitted. A value of 0 uses the system default. This parameter is
supported only on systems that support TCP_KEEPINTVL or an equivalent socket option, and on
Chapter 19. Server Configuration

Windows; on other systems, it must be zero. In sessions connected via a Unix-domain socket, this parameter is ignored and always reads as zero.

**Note:** On Windows, a value of 0 will set this parameter to 1 second, since Windows does not provide a way to read the system default value.

tcp_keepalives_count (integer)

Specifies the number of TCP keepalives that can be lost before the server’s connection to the client is considered dead. A value of 0 uses the system default. This parameter is supported only on systems that support TCP_KEEP_CNT or an equivalent socket option; on other systems, it must be zero. In sessions connected via a Unix-domain socket, this parameter is ignored and always reads as zero.

**Note:** This parameter is not supported on Windows, and must be zero.

19.3.2. Security and Authentication

authentication_timeout (integer)

Maximum time to complete client authentication, in seconds. If a would-be client has not completed the authentication protocol in this much time, the server closes the connection. This prevents hung clients from occupying a connection indefinitely. The default is one minute (1m). This parameter can only be set in the `postgresql.conf` file or on the server command line.

ssl (boolean)

Enables SSL connections. Please read Section 18.9 before using this. The default is off. This parameter can only be set at server start. SSL communication is only possible with TCP/IP connections.

ssl_ca_file (string)

Specifies the name of the file containing the SSL server certificate authority (CA). The default is empty, meaning no CA file is loaded, and client certificate verification is not performed. (In previous releases of PostgreSQL, the name of this file was hard-coded as `root.crt`.) Relative paths are relative to the data directory. This parameter can only be set at server start.

ssl_cert_file (string)

Specifies the name of the file containing the SSL server certificate. The default is `server.crt`. Relative paths are relative to the data directory. This parameter can only be set at server start.

ssl_crl_file (string)

Specifies the name of the file containing the SSL server certificate revocation list (CRL). The default is empty, meaning no CRL file is loaded. (In previous releases of PostgreSQL, the name of this file was hard-coded as `root.crl`.) Relative paths are relative to the data directory. This parameter can only be set at server start.
Chapter 19. Server Configuration

ssl_key_file (string)

Specifies the name of the file containing the SSL server private key. The default is server.key. Relative paths are relative to the data directory. This parameter can only be set at server start.

ssl_ciphers (string)

Specifies a list of SSL cipher suites that are allowed to be used on secure connections. See the ciphers manual page in the OpenSSL package for the syntax of this setting and a list of supported values. The default value is `HIGH: MEDIUM:+3DES:!aNULL`. It is usually reasonable, unless you have specific security requirements. This parameter can only be set at server start.

Explanation of the default value:

**HIGH**

Cipher suites that use ciphers from HIGH group (e.g., AES, Camellia, 3DES)

**MEDIUM**

Cipher suites that use ciphers from MEDIUM group (e.g., RC4, SEED)

**+3DES**

The OpenSSL default order for HIGH is problematic because it orders 3DES higher than AES128. This is wrong because 3DES offers less security than AES128, and it is also much slower. +3DES reorders it after all other HIGH and MEDIUM ciphers.

**!aNULL**

Disables anonymous cipher suites that do no authentication. Such cipher suites are vulnerable to man-in-the-middle attacks and therefore should not be used.

Available cipher suite details will vary across OpenSSL versions. Use the command `openssl ciphers -v 'HIGH: MEDIUM:+3DES:!aNULL'` to see actual details for the currently installed OpenSSL version. Note that this list is filtered at run time based on the server key type.

ssl_prefer_server_ciphers (bool)

Specifies whether to use the server’s SSL cipher preferences, rather than the client’s. The default is true. This parameter can only be set at server start.

Older PostgreSQL versions do not have this setting and always use the client’s preferences. This setting is mainly for backward compatibility with those versions. Using the server’s preferences is usually better because it is more likely that the server is appropriately configured.

ssl_ecdh_curve (string)

Specifies the name of the curve to use in ECDH key exchange. It needs to be supported by all clients that connect. It does not need to be same curve as used by server’s Elliptic Curve key. The default is `prime256v1`. This parameter can only be set at server start.

OpenSSL names for most common curves: `prime256v1` (NIST P-256), `secp384r1` (NIST P-384), `secp521r1` (NIST P-521).

The full list of available curves can be shown with the command `openssl ecparam -list_curves`. Not all of them are usable in TLS though.

password_encryption (boolean)

When a password is specified in CREATE USER or ALTER ROLE without writing either ENCRYPTED or UNENCRYPTED, this parameter determines whether the password is to be encrypted. The default is on (encrypt the password).
Chapter 19. Server Configuration

krb_server_keyfile (string)

Sets the location of the Kerberos server key file. See Section 20.3.3 for details. This parameter can only be set in the postgresql.conf file or on the server command line.

krb_caseins_users (boolean)

Sets whether GSSAPI user names should be treated case-insensitively. The default is off (case sensitive). This parameter can only be set in the postgresql.conf file or on the server command line.

db_user_namespace (boolean)

This parameter enables per-database user names. It is off by default. This parameter can only be set in the postgresql.conf file or on the server command line.

If this is on, you should create users as username@dbname. When username is passed by a connecting client, @ and the database name are appended to the user name and that database-specific user name is looked up by the server. Note that when you create users with names containing @ within the SQL environment, you will need to quote the user name.

With this parameter enabled, you can still create ordinary global users. Simply append @ when specifying the user name in the client, e.g. joe@. The @ will be stripped off before the user name is looked up by the server.

db_user_namespace causes the client’s and server’s user name representation to differ. Authentication checks are always done with the server’s user name so authentication methods must be configured for the server’s user name, not the client’s. Because md5 uses the user name as salt on both the client and server, md5 cannot be used with db_user_namespace.

Note: This feature is intended as a temporary measure until a complete solution is found. At that time, this option will be removed.

19.4. Resource Consumption

19.4.1. Memory

shared_buffers (integer)

Sets the amount of memory the database server uses for shared memory buffers. The default is typically 128 megabytes (128MB), but might be less if your kernel settings will not support it (as determined during initdb). This setting must be at least 128 kilobytes. (Non-default values of BLCKSZ change the minimum.) However, settings significantly higher than the minimum are usually needed for good performance. This parameter can only be set at server start.

If you have a dedicated database server with 1GB or more of RAM, a reasonable starting value for shared_buffers is 25% of the memory in your system. There are some workloads where even large settings for shared_buffers are effective, but because PostgreSQL also relies on the operating system cache, it is unlikely that an allocation of more than 40% of RAM to shared_buffers will work better than a smaller amount. Larger settings for shared_buffers
usually require a corresponding increase in `max_wal_size`, in order to spread out the process of writing large quantities of new or changed data over a longer period of time.

On systems with less than 1GB of RAM, a smaller percentage of RAM is appropriate, so as to leave adequate space for the operating system. Also, on Windows, large values for `shared_buffers` aren’t as effective. You may find better results keeping the setting relatively low and using the operating system cache more instead. The useful range for `shared_buffers` on Windows systems is generally from 64MB to 512MB.

`huge_pages (enum)`

Enables/disables the use of huge memory pages. Valid values are `try` (the default), `on`, and `off`. At present, this feature is supported only on Linux. The setting is ignored on other systems when set to `try`.

The use of huge pages results in smaller page tables and less CPU time spent on memory management, increasing performance. For more details, see Section 18.4.5.

With `huge_pages` set to `try`, the server will try to use huge pages, but fall back to using normal allocation if that fails. With `on`, failure to use huge pages will prevent the server from starting up. With `off`, huge pages will not be used.

`temp_buffers (integer)`

Sets the maximum number of temporary buffers used by each database session. These are session-local buffers used only for access to temporary tables. The default is eight megabytes (8MB). The setting can be changed within individual sessions, but only before the first use of temporary tables within the session; subsequent attempts to change the value will have no effect on that session.

A session will allocate temporary buffers as needed up to the limit given by `temp_buffers`. The cost of setting a large value in sessions that do not actually need many temporary buffers is only a buffer descriptor, or about 64 bytes, per increment in `temp_buffers`. However if a buffer is actually used an additional 8192 bytes will be consumed for it (or in general, `BLCKSZ` bytes).

`max_prepared_transactions (integer)`

Sets the maximum number of transactions that can be in the “prepared” state simultaneously (see `PREPARE TRANSACTION`). Setting this parameter to zero (which is the default) disables the prepared-transaction feature. This parameter can only be set at server start.

If you are not planning to use prepared transactions, this parameter should be set to zero to prevent accidental creation of prepared transactions. If you are using prepared transactions, you will probably want `max_prepared_transactions` to be at least as large as `max_connections`, so that every session can have a prepared transaction pending.

When running a standby server, you must set this parameter to the same or higher value than on the master server. Otherwise, queries will not be allowed in the standby server.

`work_mem (integer)`

Specifies the amount of memory to be used by internal sort operations and hash tables before writing to temporary disk files. The value defaults to four megabytes (4MB). Note that for a complex query, several sort or hash operations might be running in parallel; each operation will be allowed to use as much memory as this value specifies before it starts to write data into temporary files. Also, several running sessions could be doing such operations concurrently. Therefore, the total memory used could be many times the value of `work_mem`; it is necessary to keep this fact in mind when choosing the value. Sort operations are used for `ORDER BY`, `DISTINCT`, and merge joins. Hash tables are used in hash joins, hash-based aggregation, and hash-based processing of `IN` subqueries.
Chapter 19. Server Configuration

maintenance_work_mem(integer)

Specifies the maximum amount of memory to be used by maintenance operations, such as VACUUM, CREATE INDEX, and ALTER TABLE ADD FOREIGN KEY. It defaults to 64 megabytes (64MB). Since only one of these operations can be executed at a time by a database session, and an installation normally doesn’t have many of them running concurrently, it’s safe to set this value significantly larger than work_mem. Larger settings might improve performance for vacuuming and for restoring database dumps.

Note that when autovacuum runs, up to autovacuum_max_workers times this memory may be allocated, so be careful not to set the default value too high. It may be useful to control for this by separately setting autovacuum_work_mem.

replacement_sort_tuples(integer)

When the number of tuples to be sorted is smaller than this number, a sort will produce its first output run using replacement selection rather than quicksort. This may be useful in memory-constrained environments where tuples that are input into larger sort operations have a strong physical-to-logical correlation. Note that this does not include input tuples with an inverse correlation. It is possible for the replacement selection algorithm to generate one long run that requires no merging, where use of the default strategy would result in many runs that must be merged to produce a final sorted output. This may allow sort operations to complete sooner.

The default is 150,000 tuples. Note that higher values are typically not much more effective, and may be counter-productive, since the priority queue is sensitive to the size of available CPU cache, whereas the default strategy sorts runs using a cache oblivious algorithm. This property allows the default sort strategy to automatically and transparently make effective use of available CPU cache.

Setting maintenance_work_mem to its default value usually prevents utility command external sorts (e.g., sorts used by CREATE INDEX to build B-Tree indexes) from ever using replacement selection sort, unless the input tuples are quite wide.

autovacuum_work_mem(integer)

Specifies the maximum amount of memory to be used by each autovacuum worker process. It defaults to -1, indicating that the value of maintenance_work_mem should be used instead. The setting has no effect on the behavior of VACUUM when run in other contexts.

max_stack_depth(integer)

Specifies the maximum safe depth of the server’s execution stack. The ideal setting for this parameter is the actual stack size limit enforced by the kernel (as set by ulimit -s or local equivalent), less a safety margin of a megabyte or so. The safety margin is needed because the stack depth is not checked in every routine in the server, but only in key potentially-recursive routines such as expression evaluation. The default setting is two megabytes (2MB), which is conservatively small and unlikely to risk crashes. However, it might be too small to allow execution of complex functions. Only superusers can change this setting.

Setting max_stack_depth higher than the actual kernel limit will mean that a runaway recursive function can crash an individual backend process. On platforms where PostgreSQL can determine the kernel limit, the server will not allow this variable to be set to an unsafe value. However, not all platforms provide the information, so caution is recommended in selecting a value.

dynamic_shared_memory_type(enum)

Specifies the dynamic shared memory implementation that the server should use. Possible values are posix (for POSIX shared memory allocated using shm_open), sysv (for System V shared memory allocated via shmget), windows (for Windows shared memory), mmap (to simulate...
shared memory using memory-mapped files stored in the data directory), and none (to disable this feature). Not all values are supported on all platforms; the first supported option is the default for that platform. The use of the mmap option, which is not the default on any platform, is generally discouraged because the operating system may write modified pages back to disk repeatedly, increasing system I/O load; however, it may be useful for debugging, when the pg_dynshmem directory is stored on a RAM disk, or when other shared memory facilities are not available.

19.4.2. Disk

temp_file_limit (integer)

Specifies the maximum amount of disk space that a process can use for temporary files, such as sort and hash temporary files, or the storage file for a held cursor. A transaction attempting to exceed this limit will be canceled. The value is specified in kilobytes, and -1 (the default) means no limit. Only superusers can change this setting.

This setting constrains the total space used at any instant by all temporary files used by a given PostgreSQL process. It should be noted that disk space used for explicit temporary tables, as opposed to temporary files used behind-the-scenes in query execution, does not count against this limit.

19.4.3. Kernel Resource Usage

max_files_per_process (integer)

Sets the maximum number of simultaneously open files allowed to each server subprocess. The default is one thousand files. If the kernel is enforcing a safe per-process limit, you don’t need to worry about this setting. But on some platforms (notably, most BSD systems), the kernel will allow individual processes to open many more files than the system can actually support if many processes all try to open that many files. If you find yourself seeing “Too many open files” failures, try reducing this setting. This parameter can only be set at server start.

19.4.4. Cost-based Vacuum Delay

During the execution of VACUUM and ANALYZE commands, the system maintains an internal counter that keeps track of the estimated cost of the various I/O operations that are performed. When the accumulated cost reaches a limit (specified by vacuum_cost_limit), the process performing the operation will sleep for a short period of time, as specified by vacuum_cost_delay. Then it will reset the counter and continue execution.

The intent of this feature is to allow administrators to reduce the I/O impact of these commands on concurrent database activity. There are many situations where it is not important that maintenance commands like VACUUM and ANALYZE finish quickly; however, it is usually very important that these commands do not significantly interfere with the ability of the system to perform other database operations. Cost-based vacuum delay provides a way for administrators to achieve this.

This feature is disabled by default for manually issued VACUUM commands. To enable it, set the vacuum_cost_delay variable to a nonzero value.
vacuum_cost_delay(integer)

The length of time, in milliseconds, that the process will sleep when the cost limit has been exceeded. The default value is zero, which disables the cost-based vacuum delay feature. Positive values enable cost-based vacuuming. Note that on many systems, the effective resolution of sleep delays is 10 milliseconds; setting vacuum_cost_delay to a value that is not a multiple of 10 might have the same results as setting it to the next higher multiple of 10.

When using cost-based vacuuming, appropriate values for vacuum_cost_delay are usually quite small, perhaps 10 or 20 milliseconds. Adjusting vacuum’s resource consumption is best done by changing the other vacuum cost parameters.

vacuum_cost_page_hit(integer)

The estimated cost for vacuuming a buffer found in the shared buffer cache. It represents the cost to lock the buffer pool, lookup the shared hash table and scan the content of the page. The default value is one.

vacuum_cost_page_miss(integer)

The estimated cost for vacuuming a buffer that has to be read from disk. This represents the effort to lock the buffer pool, lookup the shared hash table, read the desired block in from the disk and scan its content. The default value is 10.

vacuum_cost_page_dirty(integer)

The estimated cost charged when vacuum modifies a block that was previously clean. It represents the extra I/O required to flush the dirty block out to disk again. The default value is 20.

vacuum_cost_limit(integer)

The accumulated cost that will cause the vacuuming process to sleep. The default value is 200.

Note: There are certain operations that hold critical locks and should therefore complete as quickly as possible. Cost-based vacuum delays do not occur during such operations. Therefore it is possible that the cost accumulates far higher than the specified limit. To avoid uselessly long delays in such cases, the actual delay is calculated as vacuum_cost_delay * accumulated_balance / vacuum_cost_limit with a maximum of vacuum_cost_delay * 4.

19.4.5. Background Writer

There is a separate server process called the background writer, whose function is to issue writes of “dirty” (new or modified) shared buffers. It writes shared buffers so server processes handling user queries seldom or never need to wait for a write to occur. However, the background writer does cause a net overall increase in I/O load, because while a repeatedly-dirtied page might otherwise be written only once per checkpoint interval, the background writer might write it several times as it is dirtied in the same interval. The parameters discussed in this subsection can be used to tune the behavior for local needs.

bgwriter_delay(integer)

Specifies the delay between activity rounds for the background writer. In each round the writer issues writes for some number of dirty buffers (controllable by the following parameters). It then sleeps for bgwriter_delay milliseconds, and repeats. When there are no dirty buffers in the buffer pool, though, it goes into a longer sleep regardless of bgwriter_delay. The default value is 200 milliseconds (200ms). Note that on many systems, the effective resolution of sleep
delays is 10 milliseconds; setting `bgwriter_delay` to a value that is not a multiple of 10 might have the same results as setting it to the next higher multiple of 10. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`bgwriter_lru_maxpages` (integer)

In each round, no more than this many buffers will be written by the background writer. Setting this to zero disables background writing. (Note that checkpoints, which are managed by a separate, dedicated auxiliary process, are unaffected.) The default value is 100 buffers. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`bgwriter_lru_multiplier` (floating point)

The number of dirty buffers written in each round is based on the number of new buffers that have been needed by server processes during recent rounds. The average recent need is multiplied by `bgwriter_lru_multiplier` to arrive at an estimate of the number of buffers that will be needed during the next round. Dirty buffers are written until there are that many clean, reusable buffers available. (However, no more than `bgwriter_lru_maxpages` buffers will be written per round.) Thus, a setting of 1.0 represents a “just in time” policy of writing exactly the number of buffers predicted to be needed. Larger values provide some cushion against spikes in demand, while smaller values intentionally leave writes to be done by server processes. The default is 2.0. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`bgwriter_flush_after` (integer)

Whenever more than `bgwriter_flush_after` bytes have been written by the bgwriter, attempt to force the OS to issue these writes to the underlying storage. Doing so will limit the amount of dirty data in the kernel’s page cache, reducing the likelihood of stalls when a fsync is issued at the end of a checkpoint, or when the OS writes data back in larger batches in the background. Often that will result in greatly reduced transaction latency, but there also are some cases, especially with workloads that are bigger than `shared_buffers`, but smaller than the OS’s page cache, where performance might degrade. This setting may have no effect on some platforms. The valid range is between 0, which disables forced writeback, and 2MB. The default is 512kB on Linux, 0 elsewhere. (If `BLCKSZ` is not 8kB, the default and maximum values scale proportionally to it.) This parameter can only be set in the `postgresql.conf` file or on the server command line.

Smaller values of `bgwriter_lru_maxpages` and `bgwriter_lru_multiplier` reduce the extra I/O load caused by the background writer, but make it more likely that server processes will have to issue writes for themselves, delaying interactive queries.

**19.4.6. Asynchronous Behavior**

`effective_io_concurrency` (integer)

Sets the number of concurrent disk I/O operations that PostgreSQL expects can be executed simultaneously. Raising this value will increase the number of I/O operations that any individual PostgreSQL session attempts to initiate in parallel. The allowed range is 1 to 1000, or zero to disable issuance of asynchronous I/O requests. Currently, this setting only affects bitmap heap scans.

For magnetic drives, a good starting point for this setting is the number of separate drives comprising a RAID 0 stripe or RAID 1 mirror being used for the database. (For RAID 5 the parity drive should not be counted.) However, if the database is often busy with multiple queries issued in concurrent sessions, lower values may be sufficient to keep the disk array busy. A value higher than needed to keep the disks busy will only result in extra CPU overhead. SSDs and other
memory-based storage can often process many concurrent requests, so the best value might be in the hundreds.

Asynchronous I/O depends on an effective `posix_fadvise` function, which some operating systems lack. If the function is not present then setting this parameter to anything but zero will result in an error. On some operating systems (e.g., Solaris), the function is present but does not actually do anything.

The default is 1 on supported systems, otherwise 0. This value can be overridden for tables in a particular tablespace by setting the tablespace parameter of the same name (see `ALTER TABLESPACE`).

### max_worker_processes (integer)

Sets the maximum number of background processes that the system can support. This parameter can only be set at server start. The default is 8.

When running a standby server, you must set this parameter to the same or higher value than on the master server. Otherwise, queries will not be allowed in the standby server.

### max_parallel_workers_per_gather (integer)

Sets the maximum number of workers that can be started by a single `Gather` node. Parallel workers are taken from the pool of processes established by `max_worker_processes`. Note that the requested number of workers may not actually be available at run time. If this occurs, the plan will run with fewer workers than expected, which may be inefficient. Setting this value to 0, which is the default, disables parallel query execution.

Note that parallel queries may consume very substantially more resources than non-parallel queries, because each worker process is a completely separate process which has roughly the same impact on the system as an additional user session. This should be taken into account when choosing a value for this setting, as well as when configuring other settings that control resource utilization, such as `work_mem`. Resource limits such as `work_mem` are applied individually to each worker, which means the total utilization may be much higher across all processes than it would normally be for any single process. For example, a parallel query using 4 workers may use up to 5 times as much CPU time, memory, I/O bandwidth, and so forth as a query which uses no workers at all.

For more information on parallel query, see Chapter 15.

### backend_flush_after (integer)

Whenever more than `backend_flush_after` bytes have been written by a single backend, attempt to force the OS to issue these writes to the underlying storage. Doing so will limit the amount of dirty data in the kernel’s page cache, reducing the likelihood of stalls when an `fsync` is issued at the end of a checkpoint, or when the OS writes data back in larger batches in the background. Often that will result in greatly reduced transaction latency, but there also are some cases, especially with workloads that are bigger than `shared_buffers`, but smaller than the OS’s page cache, where performance might degrade. This setting may have no effect on some platforms. The valid range is between 0, which disables forced writeback, and 2MB. The default is 0, i.e., no forced writeback. (If `BLCKSZ` is not 8kB, the maximum value scales proportionally to it.)

### old_snapshot_threshold (integer)

Sets the minimum time that a snapshot can be used without risk of a `snapshot too old` error occurring when using the snapshot. This parameter can only be set at server start.

Beyond the threshold, old data may be vacuumed away. This can help prevent bloat in the face of snapshots which remain in use for a long time. To prevent incorrect results due to cleanup of data which would otherwise be visible to the snapshot, an error is generated when the snapshot
is older than this threshold and the snapshot is used to read a page which has been modified since
the snapshot was built.
A value of \(-1\) disables this feature, and is the default. Useful values for production work probably
range from a small number of hours to a few days. The setting will be coerced to a granularity
of minutes, and small numbers (such as 0 or 1\text{min}) are only allowed because they may some-
times be useful for testing. While a setting as high as 60\text{d} is allowed, please note that in many
workloads extreme bloat or transaction ID wraparound may occur in much shorter time frames.
When this feature is enabled, freed space at the end of a relation cannot be released to the op-
erating system, since that could remove information needed to detect the snapshot too old
condition. All space allocated to a relation remains associated with that relation for reuse only
within that relation unless explicitly freed (for example, with \texttt{VACUUM FULL}).
This setting does not attempt to guarantee that an error will be generated under any particular
circumstances. In fact, if the correct results can be generated from (for example) a cursor which
has materialized a result set, no error will be generated even if the underlying rows in the re-
ferenced table have been vacuumed away. Some tables cannot safely be vacuumed early, and so
will not be affected by this setting. Examples include system catalogs and any table which has
a hash index. For such tables this setting will neither reduce bloat nor create a possibility of a
snapshot too old error on scanning.

19.5. Write Ahead Log

For additional information on tuning these settings, see Section 30.4.

19.5.1. Settings

\texttt{wal\_level} (enum)

\texttt{wal\_level} determines how much information is written to the WAL. The default value is
minimal, which writes only the information needed to recover from a crash or immediate shut-
down. \texttt{replica} adds logging required for WAL archiving as well as information required to run
read-only queries on a standby server. Finally, \texttt{logical} adds information necessary to support
logical decoding. Each level includes the information logged at all lower levels. This parameter
can only be set at server start.
In \texttt{minimal} level, WAL-logging of some bulk operations can be safely skipped, which can make
those operations much faster (see Section 14.4.7). Operations in which this optimization can be
applied include:

\begin{verbatim}
CREATE TABLE AS
CREATE INDEX
CLUSTER
COPY
\end{verbatim}

\texttt{COPY} into tables that were created or truncated in the same transaction

But \texttt{minimal} WAL does not contain enough information to reconstruct the data from a base
backup and the WAL logs, so \texttt{replica} or higher must be used to enable WAL archiving
(archive\_mode) and streaming replication.

In \texttt{logical} level, the same information is logged as with \texttt{replica}, plus information needed
to allow extracting logical change sets from the WAL. Using a level of \texttt{logical} will increase
the WAL volume, particularly if many tables are configured for REPLICA IDENTITY FULL and many UPDATE and DELETE statements are executed.

In releases prior to 9.6, this parameter also allowed the values archive and hot_standby. These are still accepted but mapped to replica.

**fsync (boolean)**

If this parameter is on, the PostgreSQL server will try to make sure that updates are physically written to disk, by issuing fsync() system calls or various equivalent methods (see wal_sync_method). This ensures that the database cluster can recover to a consistent state after an operating system or hardware crash.

While turning off fsync is often a performance benefit, this can result in unrecoverable data corruption in the event of a power failure or system crash. Thus it is only advisable to turn off fsync if you can easily recreate your entire database from external data.

Examples of safe circumstances for turning off fsync include the initial loading of a new database cluster from a backup file, using a database cluster for processing a batch of data after which the database will be thrown away and recreated, or for a read-only database clone which gets recreated frequently and is not used for failover. High quality hardware alone is not a sufficient justification for turning off fsync.

For reliable recovery when changing fsync off to on, it is necessary to force all modified buffers in the kernel to durable storage. This can be done while the cluster is shutdown or while fsync is on by running initdb --sync-only, running sync, unmounting the file system, or rebooting the server.

In many situations, turning off synchronous_commit for noncritical transactions can provide much of the potential performance benefit of turning off fsync, without the attendant risks of data corruption.

fsync can only be set in the postgresql.conf file or on the server command line. If you turn this parameter off, also consider turning off full_page_writes.

**synchronous_commit (enum)**

Specifies whether transaction commit will wait for WAL records to be written to disk before the command returns a “success” indication to the client. Valid values are on, remote_apply, remote_write, local, and off. The default, and safe, setting is on. When off, there can be a delay between when success is reported to the client and when the transaction is really guaranteed to be safe against a server crash. (The maximum delay is three times wal_writer_delay.) Unlike fsync, setting this parameter to off does not create any risk of database inconsistency: an operating system or database crash might result in some recent allegedly-committed transactions being lost, but the database state will be just the same as if those transactions had been aborted cleanly. So, turning synchronous_commit off can be a useful alternative when performance is more important than exact certainty about the durability of a transaction. For more discussion see Section 30.3.

If synchronous_standby_names is non-empty, this parameter also controls whether or not transaction commits will wait for their WAL records to be replicated to the standby server(s). When set to on, commits will wait until replies from the current synchronous standby(s) indicate they have received the commit record of the transaction and flushed it to disk. This ensures the transaction will not be lost unless both the primary and all synchronous standbys suffer corruption of their database storage. When set to remote_apply, commits will wait until replies from the current synchronous standby(s) indicate they have received the commit record of the transaction and applied it, so that it has become visible to queries on the standby(s). When set to remote_write, commits will wait until replies from the current synchronous standby(s) indi-
cate they have received the commit record of the transaction and written it out to their operating
system. This setting is sufficient to ensure data preservation even if a standby instance of Post-
greSQL were to crash, but not if the standby suffers an operating-system-level crash, since the
data has not necessarily reached stable storage on the standby. Finally, the setting local causes
commits to wait for local flush to disk, but not for replication. This is not usually desirable when
synchronous replication is in use, but is provided for completeness.

If synchronous_standby_names is empty, the settings on, remote_apply, remote_write
and local all provide the same synchronization level: transaction commits only wait for local
flush to disk.

This parameter can be changed at any time; the behavior for any one transaction is determined
by the setting in effect when it commits. It is therefore possible, and useful, to have some trans-
actions commit synchronously and others asynchronously. For example, to make a single multi-
statement transaction commit asynchronously when the default is the opposite, issue SET LOCAL
synchronous_commit TO OFF within the transaction.

**wal_sync_method (enum)**

Method used for forcing WAL updates out to disk. If fsync is off then this setting is irrelevant,
since WAL file updates will not be forced out at all. Possible values are:

- open_datasync (write WAL files with open() option O_DSYNC)
- fdatasync (call fdatasync() at each commit)
- fsync (call fsync() at each commit)
- fsync_writethrough (call fsync() at each commit, forcing write-through of any disk
  write cache)
- open_sync (write WAL files with open() option O_SYNC)

The open_* options also use O_DIRECT if available. Not all of these choices are available on
all platforms. The default is the first method in the above list that is supported by the platform,
except that fdatasync is the default on Linux. The default is not necessarily ideal; it might be
necessary to change this setting or other aspects of your system configuration in order to create
a crash-safe configuration or achieve optimal performance. These aspects are discussed in Section
30.1. This parameter can only be set in the postgresql.conf file or on the server command
line.

**full_page_writes (boolean)**

When this parameter is on, the PostgreSQL server writes the entire content of each disk page
to WAL during the first modification of that page after a checkpoint. This is needed because a
page write that is in process during an operating system crash might be only partially completed,
leading to an on-disk page that contains a mix of old and new data. The row-level change data
normally stored in WAL will not be enough to completely restore such a page during post-crash
recovery. Storing the full page image guarantees that the page can be correctly restored, but at
the price of increasing the amount of data that must be written to WAL. (Because WAL replay
always starts from a checkpoint, it is sufficient to do this during the first change of each page
after a checkpoint. Therefore, one way to reduce the cost of full-page writes is to increase the
checkpoint interval parameters.)

Turning this parameter off speeds normal operation, but might lead to either unrecoverable data
corruption, or silent data corruption, after a system failure. The risks are similar to turning off
fsync, though smaller, and it should be turned off only based on the same circumstances rec-
ommended for that parameter.
Turning off this parameter does not affect use of WAL archiving for point-in-time recovery (PITR) (see Section 25.3).

This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is on.

`wal_log_hints (boolean)`

When this parameter is on, the PostgreSQL server writes the entire content of each disk page to WAL during the first modification of that page after a checkpoint, even for non-critical modifications of so-called hint bits.

If data checksums are enabled, hint bit updates are always WAL-logged and this setting is ignored. You can use this setting to test how much extra WAL-logging would occur if your database had data checksums enabled.

This parameter can only be set at server start. The default value is off.

`wal_compression (boolean)`

When this parameter is on, the PostgreSQL server compresses a full page image written to WAL when `full_page_writes` is on or during a base backup. A compressed page image will be decompressed during WAL replay. The default value is off. Only superusers can change this setting.

Turning this parameter on can reduce the WAL volume without increasing the risk of unrecoverable data corruption, but at the cost of some extra CPU spent on the compression during WAL logging and on the decompression during WAL replay.

`wal_buffers (integer)`

The amount of shared memory used for WAL data that has not yet been written to disk. The default setting of -1 selects a size equal to 1/32nd (about 3%) of `shared_buffers`, but not less than 64kB nor more than the size of one WAL segment, typically 16MB. This value can be set manually if the automatic choice is too large or too small, but any positive value less than 32kB will be treated as 32kB. This parameter can only be set at server start.

The contents of the WAL buffers are written out to disk at every transaction commit, so extremely large values are unlikely to provide a significant benefit. However, setting this value to at least a few megabytes can improve write performance on a busy server where many clients are committing at once. The auto-tuning selected by the default setting of -1 should give reasonable results in most cases.

`wal_writer_delay (integer)`

Specifies how often the WAL writer flushes WAL. After flushing WAL it sleeps for `wal_writer_delay` milliseconds, unless woken up by an asynchronously committing transaction. If the last flush happened less than `wal_writer_delay` milliseconds ago and less than `wal_writer_flush_after` bytes of WAL have been produced since, then WAL is only written to the operating system, not flushed to disk. The default value is 200 milliseconds (200ms). Note that on many systems, the effective resolution of sleep delays is 10 milliseconds; setting `wal_writer_delay` to a value that is not a multiple of 10 might have the same results as setting it to the next higher multiple of 10. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`wal_writer_flush_after (integer)`

Specifies how often the WAL writer flushes WAL. If the last flush happened less than `wal_writer_delay` milliseconds ago and less than `wal_writer_flush_after` bytes of WAL have been produced since, then WAL is only written to the operating system, not flushed to disk. If `wal_writer_flush_after` is set to 0 then WAL data is flushed immediately. The
default is 1MB. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`commit_delay(integer)`

`commit_delay` adds a time delay, measured in microseconds, before a WAL flush is initiated. This can improve group commit throughput by allowing a larger number of transactions to commit via a single WAL flush, if system load is high enough that additional transactions become ready to commit within the given interval. However, it also increases latency by up to `commit_delay` microseconds for each WAL flush. Because the delay is just wasted if no other transactions become ready to commit, a delay is only performed if at least `commit_siblings` other transactions are active when a flush is about to be initiated. Also, no delays are performed if `fsync` is disabled. The default `commit_delay` is zero (no delay). Only superusers can change this setting.

In PostgreSQL releases prior to 9.3, `commit_delay` behaved differently and was much less effective: it affected only commits, rather than all WAL flushes, and waited for the entire configured delay even if the WAL flush was completed sooner. Beginning in PostgreSQL 9.3, the first process that becomes ready to flush waits for the configured interval, while subsequent processes wait only until the leader completes the flush operation.

`commit_siblings(integer)`

Minimum number of concurrent open transactions to require before performing the `commit_delay` delay. A larger value makes it more probable that at least one other transaction will become ready to commit during the delay interval. The default is five transactions.

### 19.5.2. Checkpoints

`checkpoint_timeout(integer)`

Maximum time between automatic WAL checkpoints, in seconds. The valid range is between 30 seconds and one day. The default is five minutes (5min). Increasing this parameter can increase the amount of time needed for crash recovery. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`checkpoint_completion_target(floating point)`

Specifies the target of checkpoint completion, as a fraction of total time between checkpoints. The default is 0.5. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`checkpoint_flush_after(integer)`

Whenever more than `checkpoint_flush_after` bytes have been written while performing a checkpoint, attempt to force the OS to issue these writes to the underlying storage. Doing so will limit the amount of dirty data in the kernel’s page cache, reducing the likelihood of stalls when an fsync is issued at the end of the checkpoint, or when the OS writes data back in larger batches in the background. Often that will result in greatly reduced transaction latency, but there also are some cases, especially with workloads that are bigger than `shared_buffers`, but smaller than the OS’s page cache, where performance might degrade. This setting may have no effect on some platforms. The valid range is between 0, which disables forced writeback, and 2MB. The default is 256kB on Linux, 0 elsewhere. (If `BLCKSZ` is not 8kB, the default and maximum values scale proportionally to it.) This parameter can only be set in the `postgresql.conf` file or on the server command line.
checkpoint_warning(integer)

Write a message to the server log if checkpoints caused by the filling of checkpoint segment files happen closer together than this many seconds (which suggests that max_wal_size ought to be raised). The default is 30 seconds (30s). Zero disables the warning. No warnings will be generated if checkpoint_timeout is less than checkpoint_warning. This parameter can only be set in the postgresql.conf file or on the server command line.

max_wal_size(integer)

Maximum size to let the WAL grow to between automatic WAL checkpoints. This is a soft limit; WAL size can exceed max_wal_size under special circumstances, like under heavy load, a failing archive_command, or a high wal_keep_segments setting. The default is 1 GB. Increasing this parameter can increase the amount of time needed for crash recovery. This parameter can only be set in the postgresql.conf file or on the server command line.

min_wal_size(integer)

As long as WAL disk usage stays below this setting, old WAL files are always recycled for future use at a checkpoint, rather than removed. This can be used to ensure that enough WAL space is reserved to handle spikes in WAL usage, for example when running large batch jobs. The default is 80 MB. This parameter can only be set in the postgresql.conf file or on the server command line.

19.5.3. Archiving

archive_mode(enum)

When archive_mode is enabled, completed WAL segments are sent to archive storage by setting archive_command. In addition to off, to disable, there are two modes: on, and always. During normal operation, there is no difference between the two modes, but when set to always the WAL archiver is enabled also during archive recovery or standby mode. In always mode, all files restored from the archive or streamed with streaming replication will be archived (again). See Section 26.2.9 for details.

archive_mode and archive_command are separate variables so that archive_command can be changed without leaving archiving mode. This parameter can only be set at server start. archive_mode cannot be enabled when wal_level is set to minimal.

archive_command(string)

The local shell command to execute to archive a completed WAL file segment. Any %p in the string is replaced by the path name of the file to archive, and any %f is replaced by only the file name. (The path name is relative to the working directory of the server, i.e., the cluster’s data directory.) Use % to embed an actual % character in the command. It is important for the command to return a zero exit status only if it succeeds. For more information see Section 25.3.1.

This parameter can only be set in the postgresql.conf file or on the server command line. It is ignored unless archive_mode was enabled at server start. If archive_command is an empty string (the default) while archive_mode is enabled, WAL archiving is temporarily disabled, but the server continues to accumulate WAL segment files in the expectation that a command will soon be provided. Setting archive_command to a command that does nothing but return true, e.g. /bin/true (REM on Windows), effectively disables archiving, but also breaks the chain of WAL files needed for archive recovery, so it should only be used in unusual circumstances.
archive_timeout (integer)

The archive_command is only invoked for completed WAL segments. Hence, if your server generates little WAL traffic (or has slack periods where it does so), there could be a long delay between the completion of a transaction and its safe recording in archive storage. To limit how old unarchived data can be, you can set archive_timeout to force the server to switch to a new WAL segment file periodically. When this parameter is greater than zero, the server will switch to a new segment file whenever this many seconds have elapsed since the last segment file switch, and there has been any database activity, including a single checkpoint. (Increasing checkpoint_timeout will reduce unnecessary checkpoints on an idle system.) Note that archived files that are closed early due to a forced switch are still the same length as completely full files. Therefore, it is unwise to use a very short archive_timeout — it will bloat your archive storage. archive_timeout settings of a minute or so are usually reasonable. You should consider using streaming replication, instead of archiving, if you want data to be copied off the master server more quickly than that. This parameter can only be set in the postgresql.conf file or on the server command line.

19.6. Replication

These settings control the behavior of the built-in streaming replication feature (see Section 26.2.5). Servers will be either a Master or a Standby server. Masters can send data, while Standby(s) are always receivers of replicated data. When cascading replication (see Section 26.2.7) is used, Standby server(s) can also be senders, as well as receivers. Parameters are mainly for Sending and Standby servers, though some parameters have meaning only on the Master server. Settings may vary across the cluster without problems if that is required.

19.6.1. Sending Server(s)

These parameters can be set on any server that is to send replication data to one or more standby servers. The master is always a sending server, so these parameters must always be set on the master. The role and meaning of these parameters does not change after a standby becomes the master.

max_wal_senders (integer)

Specifies the maximum number of concurrent connections from standby servers or streaming base backup clients (i.e., the maximum number of simultaneously running WAL sender processes). The default is zero, meaning replication is disabled. WAL sender processes count towards the total number of connections, so the parameter cannot be set higher than max_connections. Abrupt streaming client disconnection might cause an orphaned connection slot until a timeout is reached, so this parameter should be set slightly higher than the maximum number of expected clients so disconnected clients can immediately reconnect. This parameter can only be set at server start. wal_level must be set to replica or higher to allow connections from standby servers.

max_replication_slots (integer)

Specifies the maximum number of replication slots (see Section 26.2.6) that the server can support. The default is zero. This parameter can only be set at server start. wal_level must be set to replica or higher to allow replication slots to be used. Setting it to a lower value than the number of currently existing replication slots will prevent the server from starting.
Chapter 19. Server Configuration

wal_keep_segments (integer)

Specifies the minimum number of past log file segments kept in the pg_xlog directory, in case a standby server needs to fetch them for streaming replication. Each segment is normally 16 megabytes. If a standby server connected to the sending server falls behind by more than wal_keep_segments segments, the sending server might remove a WAL segment still needed by the standby, in which case the replication connection will be terminated. Downstream connections will also eventually fail as a result. (However, the standby server can recover by fetching the segment from archive, if WAL archiving is in use.)

This sets only the minimum number of segments retained in pg_xlog; the system might need to retain more segments for WAL archival or to recover from a checkpoint. If wal_keep_segments is zero (the default), the system doesn’t keep any extra segments for standby purposes, so the number of old WAL segments available to standby servers is a function of the location of the previous checkpoint and status of WAL archiving. This parameter can only be set in the postgresql.conf file or on the server command line.

wal_sender_timeout (integer)

Terminate replication connections that are inactive longer than the specified number of milliseconds. This is useful for the sending server to detect a standby crash or network outage. A value of zero disables the timeout mechanism. This parameter can only be set in the postgresql.conf file or on the server command line. The default value is 60 seconds.

track_commit_timestamp (bool)

Record commit time of transactions. This parameter can only be set in postgresql.conf file or on the server command line. The default value is off.

19.6.2. Master Server

These parameters can be set on the master/primary server that is to send replication data to one or more standby servers. Note that in addition to these parameters, wal_level must be set appropriately on the master server, and optionally WAL archiving can be enabled as well (see Section 19.5.3). The values of these parameters on standby servers are irrelevant, although you may wish to set them there in preparation for the possibility of a standby becoming the master.

synchronous_standby_names (string)

Specifies a list of standby servers that can support synchronous replication, as described in Section 26.2.8. There will be one or more active synchronous standbys; transactions waiting for commit will be allowed to proceed after these standby servers confirm receipt of their data. The synchronous standbys will be those whose names appear earlier in this list, and that are both currently connected and streaming data in real-time (as shown by a state of streaming in the pg_stat_replication view). Other standby servers appearing later in this list represent potential synchronous standbys. If any of the current synchronous standbys disconnects for whatever reason, it will be replaced immediately with the next-highest-priority standby. Specifying more than one standby name can allow very high availability.

This parameter specifies a list of standby servers using either of the following syntaxes:

num_sync ( standby_name [, , ] )
standby_name [, , ]

where num_sync is the number of synchronous standbys that transactions need to wait for replies from, and standby_name is the name of a standby server. For example, a setting of 3 (s1, s2, s3, s4) makes transaction commits wait until their WAL records are received by three higher-priority standbys chosen from standby servers s1, s2, s3 and s4.
The second syntax was used before PostgreSQL version 9.6 and is still supported. It’s the same as the first syntax with `num_sync` equal to 1. For example, `1 (s1, s2)` and `s1, s2` have the same meaning: either `s1` or `s2` is chosen as a synchronous standby.

The name of a standby server for this purpose is the `application_name` setting of the standby, as set in the `primary_conninfo` of the standby’s WAL receiver. There is no mechanism to enforce uniqueness. In case of duplicates one of the matching standbys will be considered as higher priority, though exactly which one is indeterminate. The special entry `*` matches any `application_name`, including the default application name of `walreceiver`.

Note: Each `standby_name` should have the form of a valid SQL identifier, unless it is `*`. You can use double-quoting if necessary. But note that `standby_names` are compared to standby application names case-insensitively, whether double-quoted or not.

If no synchronous standby names are specified here, then synchronous replication is not enabled and transaction commits will not wait for replication. This is the default configuration. Even when synchronous replication is enabled, individual transactions can be configured not to wait for replication by setting the synchronous_commit parameter to `local` or `off`.

This parameter can only be set in the `postgresql.conf` file or on the server command line.

```
vacuum_defer_cleanup_age (integer)
```

Specifies the number of transactions by which `VACUUM` and HOT updates will defer cleanup of dead row versions. The default is zero transactions, meaning that dead row versions can be removed as soon as possible, that is, as soon as they are no longer visible to any open transaction. You may wish to set this to a non-zero value on a primary server that is supporting hot standby servers, as described in Section 26.5. This allows more time for queries on the standby to complete without incurring conflicts due to early cleanup of rows. However, since the value is measured in terms of number of write transactions occurring on the primary server, it is difficult to predict just how much additional grace time will be made available to standby queries. This parameter can only be set in the `postgresql.conf` file or on the server command line.

You should also consider setting `hot_standby_feedback` on standby server(s) as an alternative to using this parameter.

This does not prevent cleanup of dead rows which have reached the age specified by `old_snapshot_threshold`.

### 19.6.3. Standby Servers

These settings control the behavior of a standby server that is to receive replication data. Their values on the master server are irrelevant.

```
hot_standby (boolean)
```

Specifies whether or not you can connect and run queries during recovery, as described in Section 26.5. The default value is `off`. This parameter can only be set at server start. It only has effect during archive recovery or in standby mode.

```
max_standby_archive_delay (integer)
```

When Hot Standby is active, this parameter determines how long the standby server should wait before canceling standby queries that conflict with about-to-be-applied WAL entries, as described in Section 26.5.2. `max_standby_archive_delay` applies when WAL data is being
read from WAL archive (and is therefore not current). The default is 30 seconds. Units are milliseconds if not specified. A value of -1 allows the standby to wait forever for conflicting queries to complete. This parameter can only be set in the `postgresql.conf` file or on the server command line.

Note that `max_standby_archive_delay` is not the same as the maximum length of time a query can run before cancellation; rather it is the maximum total time allowed to apply any one WAL segment’s data. Thus, if one query has resulted in significant delay earlier in the WAL segment, subsequent conflicting queries will have much less grace time.

`max_standby_streaming_delay (integer)`

When Hot Standby is active, this parameter determines how long the standby server should wait before canceling standby queries that conflict with about-to-be-applied WAL entries, as described in Section 26.5.2. `max_standby_streaming_delay` applies when WAL data is being received via streaming replication. The default is 30 seconds. Units are milliseconds if not specified. A value of -1 allows the standby to wait forever for conflicting queries to complete. This parameter can only be set in the `postgresql.conf` file or on the server command line.

Note that `max_standby_streaming_delay` is not the same as the maximum length of time a query can run before cancellation; rather it is the maximum total time allowed to apply WAL data once it has been received from the primary server. Thus, if one query has resulted in significant delay, subsequent conflicting queries will have much less grace time until the standby server has caught up again.

`wal_receiver_status_interval (integer)`

Specifies the minimum frequency for the WAL receiver process on the standby to send information about replication progress to the primary or upstream standby, where it can be seen using the `pg_stat_replication` view. The standby will report the last transaction log position it has written, the last position it has flushed to disk, and the last position it has applied. This parameter’s value is the maximum interval, in seconds, between reports. Updates are sent each time the write or flush positions change, or at least as often as specified by this parameter. Thus, the apply position may lag slightly behind the true position. Setting this parameter to zero disables status updates completely. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default value is 10 seconds.

`hot_standby_feedback (boolean)`

Specifies whether or not a hot standby will send feedback to the primary or upstream standby about queries currently executing on the standby. This parameter can be used to eliminate query cancels caused by cleanup records, but can cause database bloat on the primary for some workloads. Feedback messages will not be sent more frequently than once per `wal_receiver_status_interval`. The default value is `off`. This parameter can only be set in the `postgresql.conf` file or on the server command line.

If cascaded replication is in use the feedback is passed upstream until it eventually reaches the primary. Standbys make no other use of feedback they receive other than to pass upstream.

This setting does not override the behavior of `old_snapshot_threshold` on the primary; a snapshot on the standby which exceeds the primary’s age threshold can become invalid, resulting in cancellation of transactions on the standby. This is because `old_snapshot_threshold` is intended to provide an absolute limit on the time which dead rows can contribute to bloat, which would otherwise be violated because of the configuration of a standby.

`wal_receiver_timeout (integer)`

Terminate replication connections that are inactive longer than the specified number of milliseconds. This is useful for the receiving standby server to detect a primary node crash or network
outage. A value of zero disables the timeout mechanism. This parameter can only be set in the
`postgresql.conf` file or on the server command line. The default value is 60 seconds.

**wal_retrieve_retry_interval (integer)**

Specify how long the standby server should wait when WAL data is not available from any
sources (streaming replication, local `pg_xlog` or WAL archive) before retrying to retrieve WAL
data. This parameter can only be set in the `postgresql.conf` file or on the server command
line. The default value is 5 seconds. Units are milliseconds if not specified.

This parameter is useful in configurations where a node in recovery needs to control the amount
of time to wait for new WAL data to be available. For example, in archive recovery, it is possible
to make the recovery more responsive in the detection of a new WAL log file by reducing the
value of this parameter. On a system with low WAL activity, increasing it reduces the amount of
requests necessary to access WAL archives, something useful for example in cloud environments
where the amount of times an infrastructure is accessed is taken into account.

### 19.7. Query Planning

#### 19.7.1. Planner Method Configuration

These configuration parameters provide a crude method of influencing the query plans chosen by the
query optimizer. If the default plan chosen by the optimizer for a particular query is not optimal, a
temporary solution is to use one of these configuration parameters to force the optimizer to choose
a different plan. Better ways to improve the quality of the plans chosen by the optimizer include
adjusting the planer cost constants (see Section 19.7.2), running ANALYZE manually, increasing the
value of the `default_statistics_target` configuration parameter, and increasing the amount of statistics
collected for specific columns using `ALTER TABLE SET STATISTICS`.

**enable_bitmapscan (boolean)**

Enables or disables the query planner’s use of bitmap-scan plan types. The default is `on`.

**enable_hashagg (boolean)**

Enables or disables the query planner’s use of hashed aggregation plan types. The default is `on`.

**enable_hashjoin (boolean)**

Enables or disables the query planner’s use of hash-join plan types. The default is `on`.

**enable_indexscan (boolean)**

Enables or disables the query planner’s use of index-scan plan types. The default is `on`.

**enable_indexonlyscan (boolean)**

Enables or disables the query planner’s use of index-only-scan plan types (see Section 11.11).
The default is `on`.

**enable_material (boolean)**

Enables or disables the query planner’s use of materialization. It is impossible to suppress ma-
terialization entirely, but turning this variable off prevents the planner from inserting materialize
nodes except in cases where it is required for correctness. The default is `on`.

539
Chapter 19. Server Configuration

enable_mergejoin (boolean)

Enables or disables the query planner’s use of merge-join plan types. The default is on.

enable_nestloop (boolean)

Enables or disables the query planner’s use of nested-loop join plans. It is impossible to suppress nested-loop joins entirely, but turning this variable off discourages the planner from using one if there are other methods available. The default is on.

enable_seqscan (boolean)

Enables or disables the query planner’s use of sequential scan plan types. It is impossible to suppress sequential scans entirely, but turning this variable off discourages the planner from using one if there are other methods available. The default is on.

enable_sort (boolean)

Enables or disables the query planner’s use of explicit sort steps. It is impossible to suppress explicit sorts entirely, but turning this variable off discourages the planner from using one if there are other methods available. The default is on.

enable_tidscan (boolean)

Enables or disables the query planner’s use of TID scan plan types. The default is on.

19.7.2. Planner Cost Constants

The cost variables described in this section are measured on an arbitrary scale. Only their relative values matter, hence scaling them all up or down by the same factor will result in no change in the planner’s choices. By default, these cost variables are based on the cost of sequential page fetches; that is, seq_page_cost is conventionally set to 1.0 and the other cost variables are set with reference to that. But you can use a different scale if you prefer, such as actual execution times in milliseconds on a particular machine.

Note: Unfortunately, there is no well-defined method for determining ideal values for the cost variables. They are best treated as averages over the entire mix of queries that a particular installation will receive. This means that changing them on the basis of just a few experiments is very risky.

seq_page_cost (floating point)

Sets the planner’s estimate of the cost of a disk page fetch that is part of a series of sequential fetches. The default is 1.0. This value can be overridden for tables and indexes in a particular tablespace by setting the tablespace parameter of the same name (see ALTER TABLESPACE).

random_page_cost (floating point)

Sets the planner’s estimate of the cost of a non-sequentially-fetched disk page. The default is 4.0. This value can be overridden for tables and indexes in a particular tablespace by setting the tablespace parameter of the same name (see ALTER TABLESPACE).

Reducing this value relative to seq_page_cost will cause the system to prefer index scans; raising it will make index scans look relatively more expensive. You can raise or lower both values together to change the importance of disk I/O costs relative to CPU costs, which are described by the following parameters.
Random access to mechanical disk storage is normally much more expensive than four times sequential access. However, a lower default is used (4.0) because the majority of random accesses to disk, such as indexed reads, are assumed to be in cache. The default value can be thought of as modeling random access as 40 times slower than sequential, while expecting 90% of random reads to be cached.

If you believe a 90% cache rate is an incorrect assumption for your workload, you can increase random_page_cost to better reflect the true cost of random storage reads. Correspondingly, if your data is likely to be completely in cache, such as when the database is smaller than the total server memory, decreasing random_page_cost can be appropriate. Storage that has a low random read cost relative to sequential, e.g. solid-state drives, might also be better modeled with a lower value for random_page_cost.

**Tip:** Although the system will let you set random_page_cost to less than seq_page_cost, it is not physically sensible to do so. However, setting them equal makes sense if the database is entirely cached in RAM, since in that case there is no penalty for touching pages out of sequence. Also, in a heavily-cached database you should lower both values relative to the CPU parameters, since the cost of fetching a page already in RAM is much smaller than it would normally be.

cpu_tuple_cost (floating point)
Sets the planner’s estimate of the cost of processing each row during a query. The default is 0.01.

cpu_index_tuple_cost (floating point)
Sets the planner’s estimate of the cost of processing each index entry during an index scan. The default is 0.005.

cpu_operator_cost (floating point)
Sets the planner’s estimate of the cost of processing each operator or function executed during a query. The default is 0.0025.

parallel_setup_cost (floating point)
Sets the planner’s estimate of the cost of launching parallel worker processes. The default is 1000.

parallel_tuple_cost (floating point)
Sets the planner’s estimate of the cost of transferring one tuple from a parallel worker process to another process. The default is 0.1.

min_parallel_relation_size (integer)
Sets the minimum size of relations to be considered for parallel scan. The default is 8 megabytes (8MB).

effective_cache_size (integer)
Sets the planner’s assumption about the effective size of the disk cache that is available to a single query. This is factored into estimates of the cost of using an index; a higher value makes it more likely index scans will be used, a lower value makes it more likely sequential scans will be used. When setting this parameter you should consider both PostgreSQL’s shared buffers and the portion of the kernel’s disk cache that will be used for PostgreSQL data files, though some data might exist in both places. Also, take into account the expected number of concurrent queries on different tables, since they will have to share the available space. This parameter has no effect on the size of shared memory allocated by PostgreSQL, nor does it reserve kernel disk cache; it
is used only for estimation purposes. The system also does not assume data remains in the disk cache between queries. The default is 4 gigabytes (4GB).

### 19.7.3. Genetic Query Optimizer

The genetic query optimizer (GEQO) is an algorithm that does query planning using heuristic searching. This reduces planning time for complex queries (those joining many relations), at the cost of producing plans that are sometimes inferior to those found by the normal exhaustive-search algorithm. For more information see Chapter 58.

**geqo (boolean)**

Enables or disables genetic query optimization. This is on by default. It is usually best not to turn it off in production; the `geqo_threshold` variable provides more granular control of GEQO.

**geqo_threshold (integer)**

Use genetic query optimization to plan queries with at least this many `FROM` items involved. (Note that a `FULL OUTER JOIN` construct counts as only one `FROM` item.) The default is 12. For simpler queries it is usually best to use the regular, exhaustive-search planner, but for queries with many tables the exhaustive search takes too long, often longer than the penalty of executing a suboptimal plan. Thus, a threshold on the size of the query is a convenient way to manage use of GEQO.

**geqo_effort (integer)**

Controls the trade-off between planning time and query plan quality in GEQO. This variable must be an integer in the range from 1 to 10. The default value is five. Larger values increase the time spent doing query planning, but also increase the likelihood that an efficient query plan will be chosen.

`geqo_effort` doesn’t actually do anything directly; it is only used to compute the default values for the other variables that influence GEQO behavior (described below). If you prefer, you can set the other parameters by hand instead.

**geqo_pool_size (integer)**

Controls the pool size used by GEQO, that is the number of individuals in the genetic population. It must be at least two, and useful values are typically 100 to 1000. If it is set to zero (the default setting) then a suitable value is chosen based on `geqo_effort` and the number of tables in the query.

**geqo_generations (integer)**

Controls the number of generations used by GEQO, that is the number of iterations of the algorithm. It must be at least one, and useful values are in the same range as the pool size. If it is set to zero (the default setting) then a suitable value is chosen based on `geqo_pool_size`.

**geqo_selection_bias (floating point)**

Controls the selection bias used by GEQO. The selection bias is the selective pressure within the population. Values can be from 1.50 to 2.00; the latter is the default.

**geqo_seed (floating point)**

Controls the initial value of the random number generator used by GEQO to select random paths through the join order search space. The value can range from zero (the default) to one. Varying the value changes the set of join paths explored, and may result in a better or worse best path being found.
19.7.4. Other Planner Options

default_statistics_target (integer)

Sets the default statistics target for table columns without a column-specific target set via ALTER TABLE SET STATISTICS. Larger values increase the time needed to do ANALYZE, but might improve the quality of the planner’s estimates. The default is 100. For more information on the use of statistics by the PostgreSQL query planner, refer to Section 14.2.

constraint_exclusion (enum)

Controls the query planner’s use of table constraints to optimize queries. The allowed values of constraint_exclusion are on (examine constraints for all tables), off (never examine constraints), and partition (examine constraints only for inheritance child tables and UNION ALL subqueries). partition is the default setting. It is often used with inheritance and partitioned tables to improve performance.

When this parameter allows it for a particular table, the planner compares query conditions with the table’s CHECK constraints, and omits scanning tables for which the conditions contradict the constraints. For example:

CREATE TABLE parent(key integer, ...);
CREATE TABLE child1000(check (key between 1000 and 1999)) INHERITS(parent);
CREATE TABLE child2000(check (key between 2000 and 2999)) INHERITS(parent);
...
SELECT * FROM parent WHERE key = 2400;

With constraint exclusion enabled, this SELECT will not scan child1000 at all, improving performance.

Currently, constraint exclusion is enabled by default only for cases that are often used to implement table partitioning. Turning it on for all tables imposes extra planning overhead that is quite noticeable on simple queries, and most often will yield no benefit for simple queries. If you have no partitioned tables you might prefer to turn it off entirely.

Refer to Section 5.10.4 for more information on using constraint exclusion and partitioning.

cursor_tuple_fraction (floating point)

Sets the planner’s estimate of the fraction of a cursor’s rows that will be retrieved. The default is 0.1. Smaller values of this setting bias the planner towards using “fast start” plans for cursors, which will retrieve the first few rows quickly while perhaps taking a long time to fetch all rows. Larger values put more emphasis on the total estimated time. At the maximum setting of 1.0, cursors are planned exactly like regular queries, considering only the total estimated time and not how soon the first rows might be delivered.

fromCollapse_limit (integer)

The planner will merge sub-queries into upper queries if the resulting FROM list would have no more than this many items. Smaller values reduce planning time but might yield inferior query plans. The default is eight. For more information see Section 14.3.

Setting this value to geqo_threshold or more may trigger use of the GEQO planner, resulting in non-optimal plans. See Section 19.7.3.

joinCollapse_limit (integer)

The planner will rewrite explicit JOIN constructs (except FULL JOINs) into lists of FROM items whenever a list of no more than this many items would result. Smaller values reduce planning time but might yield inferior query plans.
Chapter 19. Server Configuration

By default, this variable is set the same as `from_collapse_limit`, which is appropriate for most uses. Setting it to 1 prevents any reordering of explicit JOINS. Thus, the explicit join order specified in the query will be the actual order in which the relations are joined. Because the query planner does not always choose the optimal join order, advanced users can elect to temporarily set this variable to 1, and then specify the join order they desire explicitly. For more information see Section 14.3.

Setting this value to `geqo_threshold` or more may trigger use of the GEQO planner, resulting in non-optimal plans. See Section 19.7.3.

`force_parallel_mode (enum)`

Allows the use of parallel queries for testing purposes even in cases where no performance benefit is expected. The allowed values of `force_parallel_mode` are off (use parallel mode only when it is expected to improve performance), on (force parallel query for all queries for which it is thought to be safe), and regress (like on, but with additional behavior changes as explained below).

More specifically, setting this value to `on` will add a Gather node to the top of any query plan for which this appears to be safe, so that the query runs inside of a parallel worker. Even when a parallel worker is not available or cannot be used, operations such as starting a subtransaction that would be prohibited in a parallel query context will be prohibited unless the planner believes that this will cause the query to fail. If failures or unexpected results occur when this option is set, some functions used by the query may need to be marked PARALLEL UNSAFE (or, possibly, PARALLEL RESTRICTED).

Setting this value to `regress` has all of the same effects as setting it to `on` plus some additional effects that are intended to facilitate automated regression testing. Normally, messages from a parallel worker include a context line indicating that, but a setting of `regress` suppresses this line so that the output is the same as in non-parallel execution. Also, the Gather nodes added to plans by this setting are hidden in EXPLAIN output so that the output matches what would be obtained if this setting were turned `off`.

19.8. Error Reporting and Logging

19.8.1. Where To Log

`log_destination (string)`

PostgreSQL supports several methods for logging server messages, including stderr, csvlog and syslog. On Windows, eventlog is also supported. Set this parameter to a list of desired log destinations separated by commas. The default is to log to stderr only. This parameter can only be set in the `postgresql.conf` file or on the server command line.

If csvlog is included in `log_destination`, log entries are output in “comma separated value” (CSV) format, which is convenient for loading logs into programs. See Section 19.8.4 for details. `logging_collector` must be enabled to generate CSV-format log output.

Note: On most Unix systems, you will need to alter the configuration of your system’s syslog daemon in order to make use of the syslog option for `log_destination`. PostgreSQL can log to syslog facilities `LOCAL0` through `LOCAL7` (see `syslog_facility`), but the default syslog con-
configuration on most platforms will discard all such messages. You will need to add something like:

```
local0.* /var/log/postgresql
```
to the syslog daemon's configuration file to make it work.

On Windows, when you use the `eventlog` option for `log_destination`, you should register an event source and its library with the operating system so that the Windows Event Viewer can display event log messages cleanly. See Section 18.11 for details.

**logging_collector (boolean)**

This parameter enables the logging collector, which is a background process that captures log messages sent to stderr and redirects them into log files. This approach is often more useful than logging to syslog, since some types of messages might not appear in syslog output. (One common example is dynamic-linker failure messages; another is error messages produced by scripts such as `archive_command`.) This parameter can only be set at server start.

**Note:** It is possible to log to stderr without using the logging collector; the log messages will just go to wherever the server's stderr is directed. However, that method is only suitable for low log volumes, since it provides no convenient way to rotate log files. Also, on some platforms not using the logging collector can result in lost or garbled log output, because multiple processes writing concurrently to the same log file can overwrite each other’s output.

**Note:** The logging collector is designed to never lose messages. This means that in case of extremely high load, server processes could be blocked while trying to send additional log messages when the collector has fallen behind. In contrast, syslog prefers to drop messages if it cannot write them, which means it may fail to log some messages in such cases but it will not block the rest of the system.

**log_directory (string)**

When `logging_collector` is enabled, this parameter determines the directory in which log files will be created. It can be specified as an absolute path, or relative to the cluster data directory. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is `pg_log`.

**log_filename (string)**

When `logging_collector` is enabled, this parameter sets the file names of the created log files. The value is treated as a `strftime` pattern, so `%`-escapes can be used to specify time-varying file names. (Note that if there are any time-zone-dependent `%`-escapes, the computation is done in the zone specified by `log_timezone`.) The supported `%`-escapes are similar to those listed in the Open Group’s `strftime` specification. Note that the system’s `strftime` is not used directly, so platform-specific (nonstandard) extensions do not work. The default is `postgresql-%Y-%m-%d_%H%M%S.log`.

If you specify a file name without escapes, you should plan to use a log rotation utility to avoid eventually filling the entire disk. In releases prior to 8.4, if no `%` escapes were present, PostgreSQL would append the epoch of the new log file’s creation time, but this is no longer the case.

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Chapter 19. Server Configuration

If CSV-format output is enabled in `log_destination`, `.csv` will be appended to the times-tamped log file name to create the file name for CSV-format output. (If `log_filename` ends in `.log`, the suffix is replaced instead.)

This parameter can only be set in the `postgresql.conf` file or on the server command line.

`log_file_mode` (integer)

On Unix systems this parameter sets the permissions for log files when `logging_collector` is enabled. (On Microsoft Windows this parameter is ignored.) The parameter value is expected to be a numeric mode specified in the format accepted by the `chmod` and `umask` system calls. (To use the customary octal format the number must start with a 0 (zero).)

The default permissions are 0600, meaning only the server owner can read or write the log files. The other commonly useful setting is 0640, allowing members of the owner’s group to read the files. Note however that to make use of such a setting, you’ll need to alter `log_directory` to store the files somewhere outside the cluster data directory. In any case, it’s unwise to make the log files world-readable, since they might contain sensitive data.

This parameter can only be set in the `postgresql.conf` file or on the server command line.

`log_rotation_age` (integer)

When `logging_collector` is enabled, this parameter determines the maximum lifetime of an individual log file. After this many minutes have elapsed, a new log file will be created. Set to zero to disable time-based creation of new log files. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`log_rotation_size` (integer)

When `logging_collector` is enabled, this parameter determines the maximum size of an individual log file. After this many kilobytes have been emitted into a log file, a new log file will be created. Set to zero to disable size-based creation of new log files. This parameter can only be set in the `postgresql.conf` file or on the server command line.

`log_truncate_on_rotation` (boolean)

When `logging_collector` is enabled, this parameter will cause PostgreSQL to truncate (over-write), rather than append to, any existing log file of the same name. However, truncation will occur only when a new file is being opened due to time-based rotation, not during server startup or size-based rotation. When off, pre-existing files will be appended to in all cases. For example, using this setting in combination with a `log_filename` like `postgresql-%H.log` would result in generating twenty-four hourly log files and then cyclically overwriting them. This parameter can only be set in the `postgresql.conf` file or on the server command line.

Example: To keep 7 days of logs, one log file per day named `server_log.Mon`, `server_log.Tue`, etc, and automatically overwrite last week’s log with this week’s log, set `log_filename` to `server_log.%a`, `log_truncate_on_rotation` to on, and `log_rotation_age` to 1440.

Example: To keep 24 hours of logs, one log file per hour, but also rotate sooner if the log file size exceeds 1GB, set `log_filename` to `server_log.%H%M`, `log_truncate_on_rotation` to on, `log_rotation_age` to 60, and `log_rotation_size` to 1000000. Including `%M` in `log_filename` allows any size-driven rotations that might occur to select a file name different from the hour’s initial file name.

`syslog_facility` (enum)

When logging to syslog is enabled, this parameter determines the syslog “facility” to be used. You can choose from `LOCAL0`, `LOCAL1`, `LOCAL2`, `LOCAL3`, `LOCAL4`, `LOCAL5`, `LOCAL6`, `LOCAL7`;
the default is LOCAL0. See also the documentation of your system’s syslog daemon. This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line.

\texttt{syslog\_ident (string)}

When logging to syslog is enabled, this parameter determines the program name used to identify PostgreSQL messages in syslog logs. The default is \texttt{postgres}. This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line.

\texttt{syslog\_sequence\_numbers (boolean)}

When logging to syslog and this is on (the default), then each message will be prefixed by an increasing sequence number (such as \texttt{[2]}). This circumvents the “--- last message repeated N times ---” suppression that many syslog implementations perform by default. In more modern syslog implementations, repeated message suppression can be configured (for example, \texttt{$RepeatedMsgReduction} in rsyslog), so this might not be necessary. Also, you could turn this off if you actually want to suppress repeated messages.

This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line.

\texttt{syslog\_split\_messages (boolean)}

When logging to syslog is enabled, this parameter determines how messages are delivered to syslog. When on (the default), messages are split by lines, and long lines are split so that they will fit into 1024 bytes, which is a typical size limit for traditional syslog implementations. When off, PostgreSQL server log messages are delivered to the syslog service as is, and it is up to the syslog service to cope with the potentially bulky messages.

If syslog is ultimately logging to a text file, then the effect will be the same either way, and it is best to leave the setting on, since most syslog implementations either cannot handle large messages or would need to be specially configured to handle them. But if syslog is ultimately writing into some other medium, it might be necessary or more useful to keep messages logically together.

This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line.

\texttt{event\_source (string)}

When logging to event log is enabled, this parameter determines the program name used to identify PostgreSQL messages in the log. The default is \texttt{PostgreSQL}. This parameter can only be set in the \texttt{postgresql.conf} file or on the server command line.

### 19.8.2. When To Log

\texttt{log\_min\_messages (enum)}

Controls which message levels are written to the server log. Valid values are \texttt{DEBUG5}, \texttt{DEBUG4}, \texttt{DEBUG3}, \texttt{DEBUG2}, \texttt{DEBUG1}, \texttt{INFO}, \texttt{NOTICE}, \texttt{WARNING}, \texttt{ERROR}, \texttt{LOG}, \texttt{FATAL}, and \texttt{PANIC}. Each level includes all the levels that follow it. The later the level, the fewer messages are sent to the log. The default is \texttt{WARNING}. Note that \texttt{LOG} has a different rank here than in \texttt{client\_min\_messages}. Only superusers can change this setting.

\texttt{log\_min\_error\_statement (enum)}

Controls which SQL statements that cause an error condition are recorded in the server log. The current SQL statement is included in the log entry for any message of the specified severity or higher. Valid values are \texttt{DEBUG5}, \texttt{DEBUG4}, \texttt{DEBUG3}, \texttt{DEBUG2}, \texttt{DEBUG1}, \texttt{INFO}, \texttt{NOTICE}, \texttt{WARNING}, \texttt{ERROR}, \texttt{LOG}, \texttt{FATAL}, and \texttt{PANIC}. The default is \texttt{ERROR}, which means statements causing errors,
log messages, fatal errors, or panics will be logged. To effectively turn off logging of failing statements, set this parameter to PANIC. Only superusers can change this setting.

`log_min_duration_statement (integer)`

Causes the duration of each completed statement to be logged if the statement ran for at least the specified number of milliseconds. Setting this to zero prints all statement durations. Minus-one (the default) disables logging statement durations. For example, if you set it to 250ms then all SQL statements that run 250ms or longer will be logged. Enabling this parameter can be helpful in tracking down unoptimized queries in your applications. Only superusers can change this setting.

For clients using extended query protocol, durations of the Parse, Bind, and Execute steps are logged independently.

**Note:** When using this option together with log_statement, the text of statements that are logged because of log_statement will not be repeated in the duration log message. If you are not using syslog, it is recommended that you log the PID or session ID using log_line_prefix so that you can link the statement message to the later duration message using the process ID or session ID.

Table 19-1 explains the message severity levels used by PostgreSQL. If logging output is sent to syslog or Windows’ eventlog, the severity levels are translated as shown in the table.

**Table 19-1. Message Severity Levels**

<table>
<thead>
<tr>
<th>Severity</th>
<th>Usage</th>
<th>syslog</th>
<th>eventlog</th>
</tr>
</thead>
<tbody>
<tr>
<td>DEBUG1..DEBUG5</td>
<td>Provides successively-more-detailed information for use by developers.</td>
<td>DEBUG</td>
<td>INFORMATION</td>
</tr>
<tr>
<td>INFO</td>
<td>Provides information implicitly requested by the user, e.g., output from VACUUM VERBOSE.</td>
<td>INFO</td>
<td>INFORMATION</td>
</tr>
<tr>
<td>NOTICE</td>
<td>Provides information that might be helpful to users, e.g., notice of truncation of long identifiers.</td>
<td>NOTICE</td>
<td>INFORMATION</td>
</tr>
<tr>
<td>WARNING</td>
<td>Provides warnings of likely problems, e.g., COMMIT outside a transaction block.</td>
<td>NOTICE</td>
<td>WARNING</td>
</tr>
<tr>
<td>ERROR</td>
<td>Reports an error that caused the current command to abort.</td>
<td>WARNING</td>
<td>ERROR</td>
</tr>
</tbody>
</table>
Severity | Usage | syslog | eventlog |
--- | --- | --- | --- |
LOG | Reports information of interest to administrators, e.g., checkpoint activity. | INFO | INFORMATION |
FATAL | Reports an error that caused the current session to abort. | ERR | ERROR |
PANIC | Reports an error that caused all database sessions to abort. | CRIT | ERROR |

### 19.8.3. What To Log

**application_name** *(string)*

The `application_name` can be any string of less than NAMEDATALEN characters (64 characters in a standard build). It is typically set by an application upon connection to the server. The name will be displayed in the `pg_stat_activity` view and included in CSV log entries. It can also be included in regular log entries via the `log_line_prefix` parameter. Only printable ASCII characters may be used in the `application_name` value. Other characters will be replaced with question marks (?).

**debug_print_parse** *(boolean)*

**debug_print_rewritten** *(boolean)*

**debug_print_plan** *(boolean)*

These parameters enable various debugging output to be emitted. When set, they print the resulting parse tree, the query rewriter output, or the execution plan for each executed query. These messages are emitted at LOG message level, so by default they will appear in the server log but will not be sent to the client. You can change that by adjusting `client_min_messages` and/or `log_min_messages`. These parameters are off by default.

**debug_pretty_print** *(boolean)*

When set, `debug_pretty_print` indents the messages produced by `debug_print_parse`, `debug_print_rewritten`, or `debug_print_plan`. This results in more readable but much longer output than the “compact” format used when it is off. It is on by default.

**log_checkpoints** *(boolean)*

Causes checkpoints and restartpoints to be logged in the server log. Some statistics are included in the log messages, including the number of buffers written and the time spent writing them. This parameter can only be set in the `postgresql.conf` file or on the server command line. The default is off.

**log_connections** *(boolean)*

Causes each attempted connection to the server to be logged, as well as successful completion of client authentication. Only superusers can change this parameter at session start, and it cannot be changed at all within a session. The default is off.

**Note:** Some client programs, like `psql`, attempt to connect twice while determining if a password is required, so duplicate “connection received” messages do not necessarily indicate a problem.
Chapter 19. Server Configuration

log_disconnections (boolean)

Causes session terminations to be logged. The log output provides information similar to log_connections, plus the duration of the session. Only superusers can change this parameter at session start, and it cannot be changed at all within a session. The default is off.

log_duration (boolean)

Causes the duration of every completed statement to be logged. The default is off. Only superusers can change this setting.

For clients using extended query protocol, durations of the Parse, Bind, and Execute steps are logged independently.

Note: The difference between setting this option and setting log_min_duration_statement to zero is that exceeding log_min_duration_statement forces the text of the query to be logged, but this option doesn’t. Thus, if log_duration is on and log_min_duration_statement has a positive value, all durations are logged but the query text is included only for statements exceeding the threshold. This behavior can be useful for gathering statistics in high-load installations.

log_error_verbosity (enum)

Controls the amount of detail written in the server log for each message that is logged. Valid values are TERSE, DEFAULT, and VERBOSE, each adding more fields to displayed messages. TERSE excludes the logging of DETAIL, HINT, QUERY, and CONTEXT error information. VERBOSE output includes the SQLSTATE error code (see also Appendix A) and the source code file name, function name, and line number that generated the error. Only superusers can change this setting.

log_hostname (boolean)

By default, connection log messages only show the IP address of the connecting host. Turning this parameter on causes logging of the host name as well. Note that depending on your host name resolution setup this might impose a non-negligible performance penalty. This parameter can only be set in the postgresql.conf file or on the server command line.

log_line_prefix (string)

This is a printf-style string that is output at the beginning of each log line. % characters begin “escape sequences” that are replaced with status information as outlined below. Unrecognized escapes are ignored. Other characters are copied straight to the log line. Some escapes are only recognized by session processes, and will be treated as empty by background processes such as the main server process. Status information may be aligned either left or right by specifying a numeric literal after the % and before the option. A negative value will cause the status information to be padded on the right with spaces to give it a minimum width, whereas a positive value will pad on the left. Padding can be useful to aid human readability in log files. This parameter can only be set in the postgresql.conf file or on the server command line. The default is an empty string.

<table>
<thead>
<tr>
<th>Escape</th>
<th>Effect</th>
<th>Session only</th>
</tr>
</thead>
<tbody>
<tr>
<td>%a</td>
<td>Application name</td>
<td>yes</td>
</tr>
<tr>
<td>%u</td>
<td>User name</td>
<td>yes</td>
</tr>
<tr>
<td>%d</td>
<td>Database name</td>
<td>yes</td>
</tr>
</tbody>
</table>
## Chapter 19. Server Configuration

### Escape Effect

<table>
<thead>
<tr>
<th>Escape</th>
<th>Effect</th>
<th>Session only</th>
</tr>
</thead>
<tbody>
<tr>
<td>%r</td>
<td>Remote host name or IP address, and remote port</td>
<td>yes</td>
</tr>
<tr>
<td>%h</td>
<td>Remote host name or IP address</td>
<td>yes</td>
</tr>
<tr>
<td>%p</td>
<td>Process ID</td>
<td>no</td>
</tr>
<tr>
<td>%t</td>
<td>Time stamp without milliseconds</td>
<td>no</td>
</tr>
<tr>
<td>%m</td>
<td>Time stamp with milliseconds</td>
<td>no</td>
</tr>
<tr>
<td>%n</td>
<td>Time stamp with milliseconds (as a Unix epoch)</td>
<td>no</td>
</tr>
<tr>
<td>%i</td>
<td>Command tag: type of session’s current command</td>
<td>yes</td>
</tr>
<tr>
<td>%e</td>
<td>SQLSTATE error code</td>
<td>no</td>
</tr>
<tr>
<td>%c</td>
<td>Session ID: see below</td>
<td>no</td>
</tr>
<tr>
<td>%l</td>
<td>Number of the log line for each session or process, starting at 1</td>
<td>no</td>
</tr>
<tr>
<td>%s</td>
<td>Process start time stamp</td>
<td>no</td>
</tr>
<tr>
<td>%v</td>
<td>Virtual transaction ID (backendID/localXID)</td>
<td>no</td>
</tr>
<tr>
<td>%x</td>
<td>Transaction ID (0 if none is assigned)</td>
<td>no</td>
</tr>
<tr>
<td>%q</td>
<td>Produces no output, but tells non-session processes to stop at this point in the string; ignored by session processes</td>
<td>no</td>
</tr>
<tr>
<td>%%</td>
<td>Literal %</td>
<td>no</td>
</tr>
</tbody>
</table>

The %c escape prints a quasi-unique session identifier, consisting of two 4-byte hexadecimal numbers (without leading zeros) separated by a dot. The numbers are the process start time and the process ID, so %c can also be used as a space saving way of printing those items. For example, to generate the session identifier from `pg_stat_activity`, use this query:

```sql
SELECT to_hex(trunc(EXTRACT(EPOCH FROM backend_start))::integer) || "." ||
       to_hex(pid)
FROM pg_stat_activity;
```

**Tip:** If you set a nonempty value for `log_line_prefix`, you should usually make its last character be a space, to provide visual separation from the rest of the log line. A punctuation character can be used too.

**Tip:** Syslog produces its own time stamp and process ID information, so you probably do not want to include those escapes if you are logging to syslog.
Chapter 19. Server Configuration

log_lock_waits (boolean)

Controls whether a log message is produced when a session waits longer than deadlock_timeout to acquire a lock. This is useful in determining if lock waits are causing poor performance. The default is off. Only superusers can change this setting.

log_statement (enum)

Controls which SQL statements are logged. Valid values are none (off), ddl, mod, and all (all statements). ddl logs all data definition statements, such as CREATE, ALTER, and DROP statements. mod logs all ddl statements, plus data-modifying statements such as INSERT, UPDATE, DELETE, TRUNCATE, and COPY (and PREPARE, EXECUTE, and EXPLAIN ANALYZE statements are also logged if their contained command is of an appropriate type. For clients using extended query protocol, logging occurs when an Execute message is received, and values of the Bind parameters are included (with any embedded single-quote marks doubled).

The default is none. Only superusers can change this setting.

Note: Statements that contain simple syntax errors are not logged even by the log_statement = all setting, because the log message is emitted only after basic parsing has been done to determine the statement type. In the case of extended query protocol, this setting likewise does not log statements that fail before the Execute phase (i.e., during parse analysis or planning). Set log_min_error_statement to ERROR (or lower) to log such statements.

log_replication_commands (boolean)

Causes each replication command to be logged in the server log. See Section 51.3 for more information about replication command. The default value is off. Only superusers can change this setting.

log_temp_files (integer)

Controls logging of temporary file names and sizes. Temporary files can be created for sorts, hashes, and temporary query results. A log entry is made for each temporary file when it is deleted. A value of zero logs all temporary file information, while positive values log only files whose size is greater than or equal to the specified number of kilobytes. The default setting is -1, which disables such logging. Only superusers can change this setting.

log_timezone (string)

Sets the time zone used for timestamps written in the server log. Unlike TimeZone, this value is cluster-wide, so that all sessions will report timestamps consistently. The built-in default is GMT, but that is typically overridden in postgresql.conf; initdb will install a setting there corresponding to its system environment. See Section 8.5.3 for more information. This parameter can only be set in the postgresql.conf file or on the server command line.

19.8.4. Using CSV-Format Log Output

Including csvlog in the log_destination list provides a convenient way to import log files into a database table. This option emits log lines in comma-separated-values (CSV) format, with these columns: time stamp with milliseconds, user name, database name, process ID, client host:port number, session ID, per-session line number, command tag, session start time, virtual transaction ID, regular transaction ID, error severity, SQLSTATE code, error message, error message detail, hint, internal query that led to the error (if any), character count of the error position therein, error con-
text, user query that led to the error (if any and enabled by \texttt{log_min_error_statement}), character count of the error position therein, location of the error in the PostgreSQL source code (if \texttt{log_error_verbosity} is set to \texttt{verbose}), and application name. Here is a sample table definition for storing CSV-format log output:

\begin{verbatim}
CREATE TABLE postgres_log
(
    log_time timestamp(3) with time zone,
    user_name text,
    database_name text,
    process_id integer,
    connection_from text,
    session_id text,
    session_line_num bigint,
    command_tag text,
    session_start_time timestamp with time zone,
    virtual_transaction_id text,
    transaction_id bigint,
    error_severity text,
    sql_state_code text,
    message text,
    detail text,
    hint text,
    internal_query text,
    internal_query_pos integer,
    context text,
    query text,
    query_pos integer,
    location text,
    application_name text,
    PRIMARY KEY (session_id, session_line_num)
);\end{verbatim}

To import a log file into this table, use the \texttt{COPY FROM} command:

\begin{verbatim}
COPY postgres_log FROM '/full/path/to/logfile.csv' WITH csv;
\end{verbatim}

There are a few things you need to do to simplify importing CSV log files:

1. Set \texttt{log_filename} and \texttt{log_rotation_age} to provide a consistent, predictable naming scheme for your log files. This lets you predict what the file name will be and know when an individual log file is complete and therefore ready to be imported.

2. Set \texttt{log_rotation_size} to 0 to disable size-based log rotation, as it makes the log file name difficult to predict.

3. Set \texttt{log_truncate_on_rotation} to \texttt{on} so that old log data isn’t mixed with the new in the same file.

4. The table definition above includes a primary key specification. This is useful to protect against accidentally importing the same information twice. The \texttt{COPY} command commits all of the data it imports at one time, so any error will cause the entire import to fail. If you import a partial log file and later import the file again when it is complete, the primary key violation will cause the import to fail. Wait until the log is complete and closed before importing. This procedure
Chapter 19. Server Configuration

will also protect against accidentally importing a partial line that hasn’t been completely written, which would also cause COPY to fail.

19.8.5. Process Title

These settings control how process titles of server processes are modified. Process titles are typically viewed using programs like ps or, on Windows, Process Explorer. See Section 28.1 for details.

cluster_name (string)

Sets the cluster name that appears in the process title for all server processes in this cluster. The name can be any string of less than NAMEDATALEN characters (64 characters in a standard build). Only printable ASCII characters may be used in the cluster_name value. Other characters will be replaced with question marks (?). No name is shown if this parameter is set to the empty string “” (which is the default). This parameter can only be set at server start.

update_process_title (boolean)

Enables updating of the process title every time a new SQL command is received by the server. This setting defaults to on on most platforms, but it defaults to off on Windows due to that platform’s larger overhead for updating the process title. Only superusers can change this setting.

19.9. Run-time Statistics

19.9.1. Query and Index Statistics Collector

These parameters control server-wide statistics collection features. When statistics collection is enabled, the data that is produced can be accessed via the pg_stat and pg_statio family of system views. Refer to Chapter 28 for more information.

track_activities (boolean)

Enables the collection of information on the currently executing command of each session, along with the time when that command began execution. This parameter is on by default. Note that even when enabled, this information is not visible to all users, only to superusers and the user owning the session being reported on, so it should not represent a security risk. Only superusers can change this setting.

track_activity_query_size (integer)

Specifies the number of bytes reserved to track the currently executing command for each active session, for the pg_stat_activity.query field. The default value is 1024. This parameter can only be set at server start.

track_counts (boolean)

Enables collection of statistics on database activity. This parameter is on by default, because the autovacuum daemon needs the collected information. Only superusers can change this setting.

track_io_timing (boolean)

Enables timing of database I/O calls. This parameter is off by default, because it will repeatedly query the operating system for the current time, which may cause significant overhead on some
Chapter 19. Server Configuration

Platforms. You can use the pg_test_timing tool to measure the overhead of timing on your system. I/O timing information is displayed in pg_stat_database, in the output of EXPLAIN when the Buffers option is used, and by pg_stat_statements. Only superusers can change this setting.

track_functions (enum)

Enables tracking of function call counts and time used. Specify pl to track only procedural-language functions, all to also track SQL and C language functions. The default is none, which disables function statistics tracking. Only superusers can change this setting.

Note: SQL-language functions that are simple enough to be "inlined" into the calling query will not be tracked, regardless of this setting.

stats_temp_directory (string)

Sets the directory to store temporary statistics data in. This can be a path relative to the data directory or an absolute path. The default is pg_stat_tmp. Pointing this at a RAM-based file system will decrease physical I/O requirements and can lead to improved performance. This parameter can only be set in the postgresql.conf file or on the server command line.

19.9.2. Statistics Monitoring

log_statement_stats (boolean)
log_parser_stats (boolean)
log_planner_stats (boolean)
log_executor_stats (boolean)

For each query, output performance statistics of the respective module to the server log. This is a crude profiling instrument, similar to the Unix getrusage() operating system facility. log_statement_stats reports total statement statistics, while the others report per-module statistics. log_statement_stats cannot be enabled together with any of the per-module options. All of these options are disabled by default. Only superusers can change these settings.

19.10. Automatic Vacuuming

These settings control the behavior of the autovacuum feature. Refer to Section 24.1.6 for more information. Note that many of these settings can be overridden on a per-table basis; see Storage Parameters.

autovacuum (boolean)

Controls whether the server should run the autovacuum launcher daemon. This is on by default; however, track_counts must also be enabled for autovacuum to work. This parameter can only be set in the postgresql.conf file or on the server command line; however, autovacuuming can be disabled for individual tables by changing table storage parameters.

Note that even when this parameter is disabled, the system will launch autovacuum processes if necessary to prevent transaction ID wraparound. See Section 24.1.5 for more information.
Chapter 19. Server Configuration

log_autovacuum_min_duration (integer)

Causes each action executed by autovacuum to be logged if it ran for at least the specified number of milliseconds. Setting this to zero logs all autovacuum actions. Minus-one (the default) disables logging autovacuum actions. For example, if you set this to 250ms then all automatic vacuums and analyzes that run 250ms or longer will be logged. In addition, when this parameter is set to any value other than -1, a message will be logged if an autovacuum action is skipped due to the existence of a conflicting lock. Enabling this parameter can be helpful in tracking autovacuum activity. This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_max_workers (integer)

Specifies the maximum number of autovacuum processes (other than the autovacuum launcher) that may be running at any one time. The default is three. This parameter can only be set at server start.

autovacuum_naptime (integer)

Specifies the minimum delay between autovacuum runs on any given database. In each round the daemon examines the database and issues VACUUM and ANALYZE commands as needed for tables in that database. The delay is measured in seconds, and the default is one minute (1min). This parameter can only be set in the postgresql.conf file or on the server command line.

autovacuum_vacuum_threshold (integer)

Specifies the minimum number of updated or deleted tuples needed to trigger a VACUUM in any one table. The default is 50 tuples. This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_analyze_threshold (integer)

Specifies the minimum number of inserted, updated or deleted tuples needed to trigger an ANALYZE in any one table. The default is 50 tuples. This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_vacuum_scale_factor (floating point)

Specifies a fraction of the table size to add to autovacuum_vacuum_threshold when deciding whether to trigger a VACUUM. The default is 0.2 (20% of table size). This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_analyze_scale_factor (floating point)

Specifies a fraction of the table size to add to autovacuum_analyze_threshold when deciding whether to trigger an ANALYZE. The default is 0.1 (10% of table size). This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

autovacuum_freeze_max_age (integer)

Specifies the maximum age (in transactions) that a table’s pg_class.relfrozenxid field can attain before a VACUUM operation is forced to prevent transaction ID wraparound within the table. Note that the system will launch autovacuum processes to prevent wraparound even when autovacuum is otherwise disabled.

Vacuum also allows removal of old files from the pg_clog subdirectory, which is why the default is a relatively low 200 million transactions. This parameter can only be set at server start, but the
setting can be reduced for individual tables by changing table storage parameters. For more information see Section 24.1.5.

**autovacuum_multixact_freeze_max_age (integer)**

Specifies the maximum age (in multixacts) that a table’s pg_class.relminmxid field can attain before a VACUUM operation is forced to prevent multixact ID wraparound within the table. Note that the system will launch autovacuum processes to prevent wraparound even when autovacuum is otherwise disabled.

Vacuuming multixacts also allows removal of old files from the pg_multixact/members and pg_multixact/offsets subdirectories, which is why the default is a relatively low 400 million multixacts. This parameter can only be set at server start, but the setting can be reduced for individual tables by changing table storage parameters. For more information see Section 24.1.5.1.

**autovacuum_vacuum_cost_delay (integer)**

Specifies the cost delay value that will be used in automatic VACUUM operations. If -1 is specified, the regular vacuum_cost_delay value will be used. The default value is 20 milliseconds. This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

**autovacuum_vacuum_cost_limit (integer)**

Specifies the cost limit value that will be used in automatic VACUUM operations. If -1 is specified (which is the default), the regular vacuum_cost_limit value will be used. Note that the value is distributed proportionally among the running autovacuum workers, if there is more than one, so that the sum of the limits for each worker does not exceed the value of this variable. This parameter can only be set in the postgresql.conf file or on the server command line; but the setting can be overridden for individual tables by changing table storage parameters.

## 19.11. Client Connection Defaults

### 19.11.1. Statement Behavior

**client_min_messages (enum)**

Controls which message levels are sent to the client. Valid values are DEBUG5, DEBUG4, DEBUG3, DEBUG2, DEBUG1, LOG, NOTICE, WARNING, and ERROR. Each level includes all the levels that follow it. The later the level, the fewer messages are sent. The default is NOTICE. Note that LOG has a different rank here than in log_min_messages.

INFO level messages are always sent to the client.

**search_path (string)**

This variable specifies the order in which schemas are searched when an object (table, data type, function, etc.) is referenced by a simple name with no schema specified. When there are objects of identical names in different schemas, the one found first in the search path is used. An object that is not in any of the schemas in the search path can only be referenced by specifying its containing schema with a qualified (dotted) name.

The value for search_path must be a comma-separated list of schema names. Any name that is not an existing schema, or is a schema for which the user does not have USAGE permission, is silently ignored.
If one of the list items is the special name $user, then the schema having the name returned by CURRENT_USER is substituted, if there is such a schema and the user has USAGE permission for it. (If not, $user is ignored.)

The system catalog schema, pg_catalog, is always searched, whether it is mentioned in the path or not. If it is mentioned in the path then it will be searched in the specified order. If pg_catalog is not in the path then it will be searched before searching any of the path items.

Likewise, the current session’s temporary-table schema, pg_temp_nnn, is always searched if it exists. It can be explicitly listed in the path by using the alias pg_temp. If it is not listed in the path then it is searched first (even before pg_catalog). However, the temporary schema is only searched for relation (table, view, sequence, etc) and data type names. It is never searched for function or operator names.

When objects are created without specifying a particular target schema, they will be placed in the first valid schema named in search_path. An error is reported if the search path is empty.

The default value for this parameter is "$user", public. This setting supports shared use of a database (where no users have private schemas, and all share use of public), private per-user schemas, and combinations of these. Other effects can be obtained by altering the default search path setting, either globally or per-user.

For more information on schema handling, see Section 5.8. In particular, the default configuration is suitable only when the database has a single user or a few mutually-trusting users.

The current effective value of the search path can be examined via the SQL function current_schemas (see Section 9.25). This is not quite the same as examining the value of search_path, since current_schemas shows how the items appearing in search_path were resolved.

**row_security (boolean)**

This variable controls whether to raise an error in lieu of applying a row security policy. When set to on, policies apply normally. When set to off, queries fail which would otherwise apply at least one policy. The default is on. Change to off where limited row visibility could cause incorrect results; for example, pg_dump makes that change by default. This variable has no effect on roles which bypass every row security policy, to wit, superusers and roles with the BYPASSRLS attribute.

For more information on row security policies, see CREATE POLICY.

**default_tablespace (string)**

This variable specifies the default tablespace in which to create objects (tables and indexes) when a CREATE command does not explicitly specify a tablespace.

The value is either the name of a tablespace, or an empty string to specify using the default tablespace of the current database. If the value does not match the name of any existing tablespace, PostgreSQL will automatically use the default tablespace of the current database. If a nondefault tablespace is specified, the user must have CREATE privilege for it, or creation attempts will fail.

This variable is not used for temporary tables; for them, temp_tablespaces is consulted instead.

This variable is also not used when creating databases. By default, a new database inherits its tablespace setting from the template database it is copied from.

For more information on tablespaces, see Section 22.6.

**temp_tablespaces (string)**

This variable specifies tablespaces in which to create temporary objects (temp tables and indexes on temp tables) when a CREATE command does not explicitly specify a tablespace. Temporary
files for purposes such as sorting large data sets are also created in these tablespaces.

The value is a list of names of tablespaces. When there is more than one name in the list, PostgreSQL chooses a random member of the list each time a temporary object is to be created; except that within a transaction, successively created temporary objects are placed in successive tablespaces from the list. If the selected element of the list is an empty string, PostgreSQL will automatically use the default tablespace of the current database instead.

When `temp_tablespaces` is set interactively, specifying a nonexistent tablespace is an error, as is specifying a tablespace for which the user does not have `CREATE` privilege. However, when using a previously set value, nonexistent tablespaces are ignored, as are tablespaces for which the user lacks `CREATE` privilege. In particular, this rule applies when using a value set in `postgresql.conf`.

The default value is an empty string, which results in all temporary objects being created in the default tablespace of the current database.

See also `default_tablespace`.

`check_function_bodies (boolean)`

This parameter is normally on. When set to `off`, it disables validation of the function body string during CREATE FUNCTION. Disabling validation avoids side effects of the validation process and avoids false positives due to problems such as forward references. Set this parameter to `off` before loading functions on behalf of other users; `pg_dump` does so automatically.

`default_transaction_isolation (enum)`

Each SQL transaction has an isolation level, which can be either “read uncommitted”, “read committed”, “repeatable read”, or “serializable”. This parameter controls the default isolation level of each new transaction. The default is “read committed”.

Consult Chapter 13 and SET TRANSACTION for more information.

`default_transaction_read_only (boolean)`

A read-only SQL transaction cannot alter non-temporary tables. This parameter controls the default read-only status of each new transaction. The default is `off` (read/write).

Consult SET TRANSACTION for more information.

`default_transaction_deferrable (boolean)`

When running at the `serializable` isolation level, a deferrable read-only SQL transaction may be delayed before it is allowed to proceed. However, once it begins executing it does not incur any of the overhead required to ensure serializability; so serialization code will have no reason to force it to abort because of concurrent updates, making this option suitable for long-running read-only transactions.

This parameter controls the default deferrable status of each new transaction. It currently has no effect on read-write transactions or those operating at isolation levels lower than `serializable`. The default is `off`.

Consult SET TRANSACTION for more information.

`session_replication_role (enum)`

Controls firing of replication-related triggers and rules for the current session. Setting this variable requires superuser privilege and results in discarding any previously cached query plans. Possible values are `origin` (the default), `replica` and `local`. See ALTER TABLE for more information.
statement_timeout (integer)

Abort any statement that takes more than the specified number of milliseconds, starting from the
time the command arrives at the server from the client. If log_min_error_statement is set
to ERROR or lower, the statement that timed out will also be logged. A value of zero (the default)
turns this off.

Setting statement_timeout in postgresql.conf is not recommended because it would af-
fect all sessions.

lock_timeout (integer)

Abort any statement that waits longer than the specified number of milliseconds while attempting
to acquire a lock on a table, index, row, or other database object. The time limit applies separately
to each lock acquisition attempt. The limit applies both to explicit locking requests (such as LOCK
TABLE, or SELECT FOR UPDATE without NOWAIT) and to implicitly-acquired locks. A value of
zero (the default) turns this off.

Unlike statement_timeout, this timeout can only occur while waiting for locks. Note
that if statement_timeout is nonzero, it is rather pointless to set lock_timeout
to the same or larger value, since the statement timeout would always trigger first. If
log_min_error_statement is set to ERROR or lower, the statement that timed out will be
logged.

Setting lock_timeout in postgresql.conf is not recommended because it would affect all
sessions.

idle_in_transaction_session_timeout (integer)

Terminate any session with an open transaction that has been idle for longer than the specified
duration in milliseconds. This allows any locks held by that session to be released and the con-
nection slot to be reused; it also allows tuples visible only to this transaction to be vacuumed.
See Section 24.1 for more details about this.

The default value of 0 disables this feature.

vacuum_freeze_table_age (integer)

VACUUM performs an aggressive scan if the table’s pg_class.relfrozenxid field has reached
the age specified by this setting. An aggressive scan differs from a regular VACUUM in that it
visits every page that might contain unfrozen XIDs or MXIDs, not just those that might contain
dead tuples. The default is 150 million transactions. Although users can set this value anywhere
from zero to two billions, VACUUM will silently limit the effective value to 95% of autovac-
uum_freeze_max_age, so that a periodical manual VACUUM has a chance to run before an anti-
wraparound autovacuum is launched for the table. For more information see Section 24.1.5.

vacuum_freeze_min_age (integer)

Specifies the cutoff age (in transactions) that VACUUM should use to decide whether to freeze row
versions while scanning a table. The default is 50 million transactions. Although users can set
this value anywhere from zero to one billion, VACUUM will silently limit the effective value to
half the value of autovacuum_freeze_max_age, so that there is not an unreasonably short time
between forced autovacuums. For more information see Section 24.1.5.

vacuum_multixact_freeze_table_age (integer)

VACUUM performs an aggressive scan if the table’s pg_class.relminmxid field has reached
the age specified by this setting. An aggressive scan differs from a regular VACUUM in that it
visits every page that might contain unfrozen XIDs or MXIDs, not just those that might contain
dead tuples. The default is 150 million multixacts. Although users can set this value anywhere
from zero to two billions, VACUUM will silently limit the effective value to 95% of autovac-
uum_multixact_freeze_max_age, so that a periodical manual VACUUM has a chance to run before an anti-wraparound is launched for the table. For more information see Section 24.1.5.1.

vacuum_multixact_freeze_min_age (integer)

Specifies the cutoff age (in multixacts) that VACUUM should use to decide whether to replace multixact IDs with a newer transaction ID or multixact ID while scanning a table. The default is 5 million multixacts. Although users can set this value anywhere from zero to one billion, VACUUM will silently limit the effective value to half the value of autovacuum_multixact_freeze_max_age, so that there is not an unreasonably short time between forced autovacuums. For more information see Section 24.1.5.1.

bytea_output (enum)

Sets the output format for values of type bytea. Valid values are hex (the default) and escape (the traditional PostgreSQL format). See Section 8.4 for more information. The bytea type always accepts both formats on input, regardless of this setting.

xmlbinary (enum)

Sets how binary values are to be encoded in XML. This applies for example when bytea values are converted to XML by the functions xmlelement or xmlforest. Possible values are base64 and hex, which are both defined in the XML Schema standard. The default is base64. For further information about XML-related functions, see Section 9.14.

The actual choice here is mostly a matter of taste, constrained only by possible restrictions in client applications. Both methods support all possible values, although the hex encoding will be somewhat larger than the base64 encoding.

xmloption (enum)

Sets whether DOCUMENT or CONTENT is implicit when converting between XML and character string values. See Section 8.13 for a description of this. Valid values are DOCUMENT and CONTENT. The default is CONTENT.

According to the SQL standard, the command to set this option is

```sql
SET XML OPTION { DOCUMENT | CONTENT };
```

This syntax is also available in PostgreSQL.

gin_pending_list_limit (integer)

Sets the maximum size of the GIN pending list which is used when fastupdate is enabled. If the list grows larger than this maximum size, it is cleaned up by moving the entries in it to the main GIN data structure in bulk. The default is four megabytes (4MB). This setting can be overridden for individual GIN indexes by changing index storage parameters. See Section 63.4.1 and Section 63.5 for more information.

19.11.2. Locale and Formatting

DateStyle (string)

Sets the display format for date and time values, as well as the rules for interpreting ambiguous date input values. For historical reasons, this variable contains two independent components: the output format specification (ISO, Postgres, SQL, or German) and the input/output specification for year/month/day ordering (DMY, MDY, or YMD). These can be set separately or together. The keywords Euro and European are synonyms for DMY; the keywords US, NonEuro, and NonEuropean are synonyms for MDY. See Section 8.5 for more information. The built-in default
is ISO, MDY, but initdb will initialize the configuration file with a setting that corresponds to the behavior of the chosen \texttt{lc	extunderscore time} locale.

\textbf{IntervalStyle (enum)}

Sets the display format for interval values. The value \texttt{sql	extunderscore standard} will produce output matching SQL standard interval literals. The value \texttt{postgres} (which is the default) will produce output matching PostgreSQL releases prior to 8.4 when the \texttt{DateStyle} parameter was set to \texttt{ISO}. The value \texttt{postgres	extunderscore verbose} will produce output matching PostgreSQL releases prior to 8.4 when the \texttt{DateStyle} parameter was set to non-\texttt{ISO} output. The value \texttt{iso	extunderscore 8601} will produce output matching the time interval “format with designators” defined in section 4.4.3.2 of ISO 8601.

The \texttt{IntervalStyle} parameter also affects the interpretation of ambiguous interval input. See Section 8.5.4 for more information.

\textbf{TimeZone (string)}

Sets the time zone for displaying and interpreting time stamps. The built-in default is \texttt{GMT}, but that is typically overridden in \texttt{postgresql.conf}; initdb will install a setting there corresponding to its system environment. See Section 8.5.3 for more information.

\textbf{timezone	extunderscore abbreviations (string)}

Sets the collection of time zone abbreviations that will be accepted by the server for datetime input. The default is ‘Default’, which is a collection that works in most of the world; there are also ‘Australia’ and ‘India’, and other collections can be defined for a particular installation. See Section B.4 for more information.

\textbf{extra	extunderscore float	extunderscore digits (integer)}

This parameter adjusts the number of digits displayed for floating-point values, including \texttt{float4}, \texttt{float8}, and geometric data types. The parameter value is added to the standard number of digits (\texttt{FLT	extunderscore DIG} or \texttt{DBL	extunderscore DIG} as appropriate). The value can be set as high as 3, to include partially-significant digits; this is especially useful for dumping float data that needs to be restored exactly. Or it can be set negative to suppress unwanted digits. See also Section 8.1.3.

\textbf{client	extunderscore encoding (string)}

Sets the client-side encoding (character set). The default is to use the database encoding. The character sets supported by the PostgreSQL server are described in Section 23.3.1.

\textbf{lc	extunderscore messages (string)}

Sets the language in which messages are displayed. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.

On some systems, this locale category does not exist. Setting this variable will still work, but there will be no effect. Also, there is a chance that no translated messages for the desired language exist. In that case you will continue to see the English messages.

Only superusers can change this setting, because it affects the messages sent to the server log as well as to the client, and an improper value might obscure the readability of the server logs.

\textbf{lc	extunderscore monetary (string)}

Sets the locale to use for formatting monetary amounts, for example with the \texttt{to	extunderscore char} family of functions. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.
Chapter 19. Server Configuration

lc_numeric(string)

Sets the locale to use for formatting numbers, for example with the to_char family of functions. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.

lc_time(string)

Sets the locale to use for formatting dates and times, for example with the to_char family of functions. Acceptable values are system-dependent; see Section 23.1 for more information. If this variable is set to the empty string (which is the default) then the value is inherited from the execution environment of the server in a system-dependent way.

default_text_search_config(string)

Selects the text search configuration that is used by those variants of the text search functions that do not have an explicit argument specifying the configuration. See Chapter 12 for further information. The built-in default is pg_catalog.simple, but initdb will initialize the configuration file with a setting that corresponds to the chosen lc_ctype locale, if a configuration matching that locale can be identified.

19.11.3. Shared Library Preloading

Several settings are available for preloading shared libraries into the server, in order to load additional functionality or achieve performance benefits. For example, a setting of \'$libdir/mylib\' would cause mylib.so (or on some platforms, mylib.sl) to be preloaded from the installation’s standard library directory. The differences between the settings are when they take effect and what privileges are required to change them.

PostgreSQL procedural language libraries can be preloaded in this way, typically by using the syntax \'$libdir/plXXX\' where XXX is pgsql, perl, tcl, or python.

For each parameter, if more than one library is to be loaded, separate their names with commas. All library names are converted to lower case unless double-quoted.

Only shared libraries specifically intended to be used with PostgreSQL can be loaded this way. Every PostgreSQL-supported library has a “magic block” that is checked to guarantee compatibility. For this reason, non-PostgreSQL libraries cannot be loaded in this way. You might be able to use operating-system facilities such as LD_PRELOAD for that.

In general, refer to the documentation of a specific module for the recommended way to load that module.

local_preload_libraries(string)

This variable specifies one or more shared libraries that are to be preloaded at connection start. The parameter value only takes effect at the start of the connection. Subsequent changes have no effect. If a specified library is not found, the connection attempt will fail.

This option can be set by any user. Because of that, the libraries that can be loaded are restricted to those appearing in the plugins subdirectory of the installation’s standard library directory. (It is the database administrator’s responsibility to ensure that only “safe” libraries are installed there.) Entries in local_preload_libraries can specify this directory explicitly, for example \$libdir/plugins/mylib, or just specify the library name — mylib would have the same effect as \$libdir/plugins/mylib.
The intent of this feature is to allow unprivileged users to load debugging or performance-measurement libraries into specific sessions without requiring an explicit LOAD command. To that end, it would be typical to set this parameter using the PGOPTIONS environment variable on the client or by using ALTER ROLE SET.

However, unless a module is specifically designed to be used in this way by non-superusers, this is usually not the right setting to use. Look at session_preload_libraries instead.

**session_preload_libraries** (string)

This variable specifies one or more shared libraries that are to be preloaded at connection start. Only superusers can change this setting. The parameter value only takes effect at the start of the connection. Subsequent changes have no effect. If a specified library is not found, the connection attempt will fail.

The intent of this feature is to allow debugging or performance-measurement libraries to be loaded into specific sessions without an explicit LOAD command being given. For example, auto_explain could be enabled for all sessions under a given user name by setting this parameter with ALTER ROLE SET. Also, this parameter can be changed without restarting the server (but changes only take effect when a new session is started), so it is easier to add new modules this way, even if they should apply to all sessions.

Unlike shared_preload_libraries, there is no large performance advantage to loading a library at session start rather than when it is first used. There is some advantage, however, when connection pooling is used.

**shared_preload_libraries** (string)

This variable specifies one or more shared libraries to be preloaded at server start. This parameter can only be set at server start. If a specified library is not found, the server will fail to start.

Some libraries need to perform certain operations that can only take place at postmaster start, such as allocating shared memory, reserving light-weight locks, or starting background workers. Those libraries must be loaded at server start through this parameter. See the documentation of each library for details.

Other libraries can also be preloaded. By preloading a shared library, the library startup time is avoided when the library is first used. However, the time to start each new server process might increase slightly, even if that process never uses the library. So this parameter is recommended only for libraries that will be used in most sessions. Also, changing this parameter requires a server restart, so this is not the right setting to use for short-term debugging tasks, say. Use session_preload_libraries for that instead.

**Note:** On Windows hosts, preloading a library at server start will not reduce the time required to start each new server process; each server process will re-load all preload libraries. However, shared_preload_libraries is still useful on Windows hosts for libraries that need to perform operations at postmaster start time.

## 19.11.4. Other Defaults
Chapter 19. Server Configuration

dynamic_library_path (string)

If a dynamically loadable module needs to be opened and the file name specified in the **CREATE FUNCTION** or **LOAD** command does not have a directory component (i.e., the name does not contain a slash), the system will search this path for the required file.

The value for **dynamic_library_path** must be a list of absolute directory paths separated by colons (or semi-colons on Windows). If a list element starts with the special string `$libdir`, the compiled-in PostgreSQL package library directory is substituted for `$libdir`; this is where the modules provided by the standard PostgreSQL distribution are installed. (Use `pg_config --pkglibdir` to find out the name of this directory.) For example:

```
dynamic_library_path = '/usr/local/lib/postgresql:/home/my_project/lib:$libdir'
```

or, in a Windows environment:

```
dynamic_library_path = 'C:\tools\postgresql;H:\my_project\lib;$libdir'
```

The default value for this parameter is `$libdir`. If the value is set to an empty string, the automatic path search is turned off.

This parameter can be changed at run time by superusers, but a setting done that way will only persist until the end of the client connection, so this method should be reserved for development purposes. The recommended way to set this parameter is in the `postgresql.conf` configuration file.

gin_fuzzy_search_limit (integer)

Soft upper limit of the size of the set returned by GIN index scans. For more information see Section 63.5.

19.12. Lock Management

deadlock_timeout (integer)

This is the amount of time, in milliseconds, to wait on a lock before checking to see if there is a deadlock condition. The check for deadlock is relatively expensive, so the server doesn’t run it every time it waits for a lock. We optimistically assume that deadlocks are not common in production applications and just wait on the lock for a while before checking for a deadlock. Increasing this value reduces the amount of time wasted in needless deadlock checks, but slows down reporting of real deadlock errors. The default is one second (1s), which is probably about the smallest value you would want in practice. On a heavily loaded server you might want to raise it. Ideally the setting should exceed your typical transaction time, so as to improve the odds that a lock will be released before the waiter decides to check for deadlock. Only superusers can change this setting.

When `log_lock_waits` is set, this parameter also determines the length of time to wait before a log message is issued about the lock wait. If you are trying to investigate locking delays you might want to set a shorter than normal `deadlock_timeout`.

max_locks_per_transaction (integer)

The shared lock table tracks locks on `max_locks_per_transaction * (max_connections + max_prepared_transactions)` objects (e.g., tables); hence, no more than this many distinct objects can be locked at any one time. This parameter controls the average number of object locks allocated for each transaction; individual transactions can lock more objects as long as the locks of all transactions fit in the lock table. This is not the number of rows that can be locked; that
value is unlimited. The default, 64, has historically proven sufficient, but you might need to raise this value if you have queries that touch many different tables in a single transaction, e.g. query of a parent table with many children. This parameter can only be set at server start.

When running a standby server, you must set this parameter to the same or higher value than on the master server. Otherwise, queries will not be allowed in the standby server.

**max_pred_locks_per_transaction (integer)**

The shared predicate lock table tracks locks on \(\text{max_pred_locks_per_transaction} \times (\text{max_connections} + \text{max_prepared_transactions})\) objects (e.g., tables); hence, no more than this many distinct objects can be locked at any one time. This parameter controls the average number of object locks allocated for each transaction; individual transactions can lock more objects as long as the locks of all transactions fit in the lock table. This is not the number of rows that can be locked; that value is unlimited. The default, 64, has generally been sufficient in testing, but you might need to raise this value if you have clients that touch many different tables in a single serializable transaction. This parameter can only be set at server start.

### 19.13. Version and Platform Compatibility

#### 19.13.1. Previous PostgreSQL Versions

**array_nulls (boolean)**

This controls whether the array input parser recognizes unquoted NULL as specifying a null array element. By default, this is on, allowing array values containing null values to be entered. However, PostgreSQL versions before 8.2 did not support null values in arrays, and therefore would treat NULL as specifying a normal array element with the string value “NULL”. For backward compatibility with applications that require the old behavior, this variable can be turned off.

Note that it is possible to create array values containing null values even when this variable is off.

**backslash_quote (enum)**

This controls whether a quote mark can be represented by \' in a string literal. The preferred, SQL-standard way to represent a quote mark is by doubling it (\") but PostgreSQL has historically also accepted \'. However, use of \' creates security risks because in some client character set encodings, there are multibyte characters in which the last byte is numerically equivalent to ASCII \. If client-side code does escaping incorrectly then a SQL-injection attack is possible. This risk can be prevented by making the server reject queries in which a quote mark appears to be escaped by a backslash. The allowed values of backslash_quote are on (allow \' always), off (reject always), and safe_encoding (allow only if client encoding does not allow ASCII \\ within a multibyte character). safe_encoding is the default setting.

Note that in a standard-conforming string literal, \ just means \\ anyway. This parameter only affects the handling of non-standard-conforming literals, including escape string syntax (E’ . .’).

**default_with_oids (boolean)**

This controls whether CREATE TABLE and CREATE TABLE AS include an OID column in newly-created tables, if neither WITH OIDS nor WITHOUT OIDS is specified. It also determines whether OIDs will be included in tables created by SELECT INTO. The parameter is off by default; in PostgreSQL 8.0 and earlier, it was on by default.
The use of OIDs in user tables is considered deprecated, so most installations should leave this variable disabled. Applications that require OIDs for a particular table should specify `WITH OIDS` when creating the table. This variable can be enabled for compatibility with old applications that do not follow this behavior.

**escape_string_warning(boolean)**

When on, a warning is issued if a backslash (\) appears in an ordinary string literal (‘...’ syntax) and `standard_conforming_strings` is off. The default is on.

Applications that wish to use backslash as escape should be modified to use escape string syntax (E’...’), because the default behavior of ordinary strings is now to treat backslash as an ordinary character, per SQL standard. This variable can be enabled to help locate code that needs to be changed.

**lo_compat_privileges(boolean)**

In PostgreSQL releases prior to 9.0, large objects did not have access privileges and were, therefore, always readable and writable by all users. Setting this variable to on disables the new privilege checks, for compatibility with prior releases. The default is off. Only superusers can change this setting.

Setting this variable does not disable all security checks related to large objects — only those for which the default behavior has changed in PostgreSQL 9.0. For example, `lo_import()` and `lo_export()` need superuser privileges regardless of this setting.

**operator_precedence_warning(boolean)**

When on, the parser will emit a warning for any construct that might have changed meanings since PostgreSQL 9.4 as a result of changes in operator precedence. This is useful for auditing applications to see if precedence changes have broken anything; but it is not meant to be kept turned on in production, since it will warn about some perfectly valid, standard-compliant SQL code. The default is off.

See Section 4.1.6 for more information.

**quote_all_identifiers(boolean)**

When the database generates SQL, force all identifiers to be quoted, even if they are not (currently) keywords. This will affect the output of `EXPLAIN` as well as the results of functions like `pg_get_viewdef`. See also the `--quote-all-identifiers` option of `pg_dump` and `pg_dumpall`.

**sql_inheritance(boolean)**

This setting controls whether undecorated table references are considered to include inheritance child tables. The default is on, which means child tables are included (thus, a * suffix is assumed by default). If turned off, child tables are not included (thus, an ONLY prefix is assumed). The SQL standard requires child tables to be included, so the off setting is not spec-compliant, but it is provided for compatibility with PostgreSQL releases prior to 7.1. See Section 5.9 for more information.

Turning `sql_inheritance` off is deprecated, because that behavior has been found to be error-prone as well as contrary to SQL standard. Discussions of inheritance behavior elsewhere in this manual generally assume that it is on.

**standard_conforming_strings(boolean)**

This controls whether ordinary string literals (‘...’) treat backslashes literally, as specified in the SQL standard. Beginning in PostgreSQL 9.1, the default is on (prior releases defaulted to off). Applications can check this parameter to determine how string literals will be processed.
Chapter 19. Server Configuration

The presence of this parameter can also be taken as an indication that the escape string syntax (`E'...'`) is supported. Escape string syntax (Section 4.1.2.2) should be used if an application desires backslashes to be treated as escape characters.

`synchronize_seqscans (boolean)`

This allows sequential scans of large tables to synchronize with each other, so that concurrent scans read the same block at about the same time and hence share the I/O workload. When this is enabled, a scan might start in the middle of the table and then “wrap around” the end to cover all rows, so as to synchronize with the activity of scans already in progress. This can result in unpredictable changes in the row ordering returned by queries that have no ORDER BY clause. Setting this parameter to off ensures the pre-8.3 behavior in which a sequential scan always starts from the beginning of the table. The default is on.

### 19.13.2. Platform and Client Compatibility

`transform_null_equals (boolean)`

When on, expressions of the form `expr = NULL` (or `NULL = expr`) are treated as `expr IS NULL`, that is, they return true if `expr` evaluates to the null value, and false otherwise. The correct SQL-spec-compliant behavior of `expr = NULL` is to always return null (unknown). Therefore this parameter defaults to off.

However, filtered forms in Microsoft Access generate queries that appear to use `expr = NULL` to test for null values, so if you use that interface to access the database you might want to turn this option on. Since expressions of the form `expr = NULL` always return the null value (using the SQL standard interpretation), they are not very useful and do not appear often in normal applications so this option does little harm in practice. But new users are frequently confused about the semantics of expressions involving null values, so this option is off by default.

Note that this option only affects the exact form `= NULL`, not other comparison operators or other expressions that are computationally equivalent to some expression involving the equals operator (such as `IN`). Thus, this option is not a general fix for bad programming.

Refer to Section 9.2 for related information.

### 19.14. Error Handling

`exit_on_error (boolean)`

If true, any error will terminate the current session. By default, this is set to false, so that only FATAL errors will terminate the session.

`restart_after_crash (boolean)`

When set to true, which is the default, PostgreSQL will automatically reinitialize after a backend crash. Leaving this value set to true is normally the best way to maximize the availability of the database. However, in some circumstances, such as when PostgreSQL is being invoked by clusterware, it may be useful to disable the restart so that the clusterware can gain control and take any actions it deems appropriate.
Chapter 19. Server Configuration

data_sync_retry (boolean)

When set to false, which is the default, PostgreSQL will raise a PANIC-level error on failure to flush modified data files to the filesystem. This causes the database server to crash. This parameter can only be set at server start.

On some operating systems, the status of data in the kernel’s page cache is unknown after a write-back failure. In some cases it might have been entirely forgotten, making it unsafe to retry; the second attempt may be reported as successful, when in fact the data has been lost. In these circumstances, the only way to avoid data loss is to recover from the WAL after any failure is reported, preferably after investigating the root cause of the failure and replacing any faulty hardware.

If set to true, PostgreSQL will instead report an error but continue to run so that the data flushing operation can be retried in a later checkpoint. Only set it to true after investigating the operating system’s treatment of buffered data in case of write-back failure.

19.15. Preset Options

The following “parameters” are read-only, and are determined when PostgreSQL is compiled or when it is installed. As such, they have been excluded from the sample `postgresql.conf` file. These options report various aspects of PostgreSQL behavior that might be of interest to certain applications, particularly administrative front-ends.

block_size (integer)

Reports the size of a disk block. It is determined by the value of `BLCKSZ` when building the server. The default value is 8192 bytes. The meaning of some configuration variables (such as `shared_buffers`) is influenced by `block_size`. See Section 19.4 for information.

data_checksums (boolean)

Reports whether data checksums are enabled for this cluster. See data checksums for more information.

debuge.assertions (boolean)

Reports whether PostgreSQL has been built with assertions enabled. That is the case if the macro `USE_ASSERT_CHECKING` is defined when PostgreSQL is built (accomplished e.g. by the `configure` option `--enable-cassert`). By default PostgreSQL is built without assertions.

integer_datetimes (boolean)

Reports whether PostgreSQL was built with support for 64-bit-integer dates and times. This can be disabled by configuring with `--disable-integer-datetimes` when building PostgreSQL. The default value is on.

lc_collate (string)

Reports the locale in which sorting of textual data is done. See Section 23.1 for more information. This value is determined when a database is created.

lc_ctype (string)

Reports the locale that determines character classifications. See Section 23.1 for more information. This value is determined when a database is created. Ordinarily this will be the same as `lc_collate`, but for special applications it might be set differently.
max_function_args (integer)

Reports the maximum number of function arguments. It is determined by the value of
FUNC_MAX_ARGS when building the server. The default value is 100 arguments.

max_identifier_length (integer)

Reports the maximum identifier length. It is determined as one less than the value of
NAMEDATALEN when building the server. The default value of NAMEDATALEN is 64; therefore
the default max_identifier_length is 63 bytes, which can be less than 63 characters when
using multibyte encodings.

max_index_keys (integer)

Reports the maximum number of index keys. It is determined by the value of INDEX_MAX_KEYS
when building the server. The default value is 32 keys.

segment_size (integer)

Reports the number of blocks (pages) that can be stored within a file segment. It is determined
by the value of RELSEG_SIZE when building the server. The maximum size of a segment file in
bytes is equal to segment_size multiplied by block_size; by default this is 1GB.

server_encoding (string)

Reports the database encoding (character set). It is determined when the database is created.
Ordinarily, clients need only be concerned with the value of client_encoding.

server_version (string)

Reports the version number of the server. It is determined by the value of PG_VERSION when
building the server.

server_version_num (integer)

Reports the version number of the server as an integer. It is determined by the value of
PG_VERSION_NUM when building the server.

wal_block_size (integer)

Reports the size of a WAL disk block. It is determined by the value of XLOG_BLCKSZ when
building the server. The default value is 8192 bytes.

wal_segment_size (integer)

Reports the number of blocks (pages) in a WAL segment file. The total size of a WAL segment
file in bytes is equal to wal_segment_size multiplied by wal_block_size; by default this is
16MB. See Section 30.4 for more information.

19.16. Customized Options

This feature was designed to allow parameters not normally known to PostgreSQL to be added by
add-on modules (such as procedural languages). This allows extension modules to be configured in
the standard ways.

Custom options have two-part names: an extension name, then a dot, then the parameter name proper,
much like qualified names in SQL. An example is plpgsql.variable_conflict.

Because custom options may need to be set in processes that have not loaded the relevant extension
module, PostgreSQL will accept a setting for any two-part parameter name. Such variables are treated
as placeholders and have no function until the module that defines them is loaded. When an extension
module is loaded, it will add its variable definitions, convert any placeholder values according to those
definitions, and issue warnings for any unrecognized placeholders that begin with its extension name.

19.17. Developer Options

The following parameters are intended for work on the PostgreSQL source code, and in some cases
to assist with recovery of severely damaged databases. There should be no reason to use them on
a production database. As such, they have been excluded from the sample postgresql.conf file.
Note that many of these parameters require special source compilation flags to work at all.

allow_system_table_mods (boolean)

Allows modification of the structure of system tables. This is used by initdb. This parameter
can only be set at server start.

ignore_system_indexes (boolean)

Ignore system indexes when reading system tables (but still update the indexes when modifying
the tables). This is useful when recovering from damaged system indexes. This parameter cannot
be changed after session start.

post_auth_delay (integer)

If nonzero, a delay of this many seconds occurs when a new server process is started, after it
conducts the authentication procedure. This is intended to give developers an opportunity to
attach to the server process with a debugger. This parameter cannot be changed after session start.

pre_auth_delay (integer)

If nonzero, a delay of this many seconds occurs just after a new server process is forked, before
it conducts the authentication procedure. This is intended to give developers an opportunity to
attach to the server process with a debugger to trace down misbehavior in authentication. This
parameter can only be set in the postgresql.conf file or on the server command line.

trace_notify (boolean)

Generates a great amount of debugging output for the LISTEN and NOTIFY commands.
client_min_messages or log_min_messages must be DEBUG1 or lower to send this output to the
client or server logs, respectively.

trace_recovery_messages (enum)

Enables logging of recovery-related debugging output that otherwise would not be logged. This
parameter allows the user to override the normal setting of log_min_messages, but only for spe-
cific messages. This is intended for use in debugging Hot Standby. Valid values are DEBUG5,
DEBUG4, DEBUG3, DEBUG2, DEBUG1, and LOG. The default, LOG, does not affect logging deci-
sions at all. The other values cause recovery-related debug messages of that priority or higher
to be logged as though they had LOG priority; for common settings of log_min_messages this
results in unconditionally sending them to the server log. This parameter can only be set in the
postgresql.conf file or on the server command line.

trace_sort (boolean)

If on, emit information about resource usage during sort operations. This parameter is only
available if the TRACE_SORT macro was defined when PostgreSQL was compiled. (However,
TRACE_SORT is currently defined by default.)
trace_locks (boolean)

If on, emit information about lock usage. Information dumped includes the type of lock operation, the type of lock and the unique identifier of the object being locked or unlocked. Also included are bit masks for the lock types already granted on this object as well as for the lock types awaited on this object. For each lock type a count of the number of granted locks and waiting locks is also dumped as well as the totals. An example of the log file output is shown here:

LOG: LockAcquire: new: lock(0xb7acd844) id(24688,24696,0,0,0,1)
     grantMask(0) req(0,0,0,0,0,0,0) grant(0,0,0,0,0,0,0)=0
     wait(0) type(AccessShareLock)

LOG: GrantLock: lock(0xb7acd844) id(24688,24696,0,0,0,1)
     grantMask(2) req(1,0,0,0,0,0,0)=1 grant(1,0,0,0,0,0,0)=1
     wait(0) type(AccessShareLock)

LOG: UnGrantLock: updated: lock(0xb7acd844) id(24688,24696,0,0,0,1)
     grantMask(0) req(0,0,0,0,0,0,0)=0 grant(0,0,0,0,0,0,0)=0
     wait(0) type(AccessShareLock)

LOG: CleanUpLock: deleting: lock(0xb7acd844) id(24688,24696,0,0,0,1)
     grantMask(0) req(0,0,0,0,0,0,0)=0 grant(0,0,0,0,0,0,0)=0
     wait(0) type(INVALID)

Details of the structure being dumped may be found in src/include/storage/lock.h.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

trace_lwlocks (boolean)

If on, emit information about lightweight lock usage. Lightweight locks are intended primarily to provide mutual exclusion of access to shared-memory data structures.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

trace_userlocks (boolean)

If on, emit information about user lock usage. Output is the same as for trace_locks, only for advisory locks.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

trace_lock_oidmin (integer)

If set, do not trace locks for tables below this OID. (use to avoid output on system tables)

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

trace_lock_table (integer)

Unconditionally trace locks on this table (OID).

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.

dump_deadlocks (boolean)

If set, dumps information about all current locks when a deadlock timeout occurs.

This parameter is only available if the LOCK_DEBUG macro was defined when PostgreSQL was compiled.
log_btree_build_stats (boolean)

If set, logs system resource usage statistics (memory and CPU) on various B-tree operations.

This parameter is only available if the `BTREE_BUILD_STATS` macro was defined when PostgreSQL was compiled.

wal_debug (boolean)

If on, emit WAL-related debugging output. This parameter is only available if the `WAL_DEBUG` macro was defined when PostgreSQL was compiled.

ignore_checksum_failure (boolean)

Only has effect if data checksums are enabled.

Detection of a checksum failure during a read normally causes PostgreSQL to report an error, aborting the current transaction. Setting `ignore_checksum_failure` to on causes the system to ignore the failure (but still report a warning), and continue processing. This behavior may cause crashes, propagate or hide corruption, or other serious problems. However, it may allow you to get past the error and retrieve undamaged tuples that might still be present in the table if the block header is still sane. If the header is corrupt an error will be reported even if this option is enabled. The default setting is `off`, and it can only be changed by a superuser.

zero_damaged_pages (boolean)

Detection of a damaged page header normally causes PostgreSQL to report an error, aborting the current transaction. Setting `zero_damaged_pages` to on causes the system to instead report a warning, zero out the damaged page in memory, and continue processing. This behavior will destroy data, namely all the rows on the damaged page. However, it does allow you to get past the error and retrieve rows from any undamaged pages that might be present in the table. It is useful for recovering data if corruption has occurred due to a hardware or software error. You should generally not set this on until you have given up hope of recovering data from the damaged pages of a table. Zeroed-out pages are not forced to disk so it is recommended to recreate the table or the index before turning this parameter off again. The default setting is `off`, and it can only be changed by a superuser.

### 19.18. Short Options

For convenience there are also single letter command-line option switches available for some parameters. They are described in Table 19-2. Some of these options exist for historical reasons, and their presence as a single-letter option does not necessarily indicate an endorsement to use the option heavily.

<table>
<thead>
<tr>
<th>Short Option</th>
<th>Equivalent</th>
</tr>
</thead>
<tbody>
<tr>
<td>-B x</td>
<td><code>shared_buffers = x</code></td>
</tr>
<tr>
<td>-d x</td>
<td><code>log_min_messages = DEBUGx</code></td>
</tr>
<tr>
<td>-e</td>
<td><code>datestyle = euro</code></td>
</tr>
</tbody>
</table>
## Chapter 19. Server Configuration

<table>
<thead>
<tr>
<th>Short Option</th>
<th>Equivalent</th>
</tr>
</thead>
<tbody>
<tr>
<td>-fb, -fh, -fi, -fm, -fn, -fo, -fs, -ft</td>
<td>enable_bitmapscan = off, enable_hashjoin = off, enable_indexscan = off, enable_mergejoin = off, enable_nestloop = off, enable_indexonlyscan = off, enable_seqscan = off, enable_tidscan = off</td>
</tr>
<tr>
<td>-F</td>
<td>fsync = off</td>
</tr>
<tr>
<td>-h x</td>
<td>listen_addresses = x</td>
</tr>
<tr>
<td>-i</td>
<td>listen_addresses = '*'</td>
</tr>
<tr>
<td>-k x</td>
<td>unix_socket_directories = x</td>
</tr>
<tr>
<td>-l</td>
<td>ssl = on</td>
</tr>
<tr>
<td>-N x</td>
<td>max_connections = x</td>
</tr>
<tr>
<td>-O</td>
<td>allow_system_table_mods = on</td>
</tr>
<tr>
<td>-p x</td>
<td>port = x</td>
</tr>
<tr>
<td>-P</td>
<td>ignore_system_indexes = on</td>
</tr>
<tr>
<td>-s</td>
<td>log_statement_stats = on</td>
</tr>
<tr>
<td>-S x</td>
<td>work_mem = x</td>
</tr>
<tr>
<td>-tpa, -tpl, -te</td>
<td>log_parser_stats = on, log_planner_stats = on, log_executor_stats = on</td>
</tr>
<tr>
<td>-W x</td>
<td>post_auth_delay = x</td>
</tr>
</tbody>
</table>
Chapter 20. Client Authentication

When a client application connects to the database server, it specifies which PostgreSQL database user name it wants to connect as, much the same way one logs into a Unix computer as a particular user. Within the SQL environment the active database user name determines access privileges to database objects — see Chapter 21 for more information. Therefore, it is essential to restrict which database users can connect.

**Note:** As explained in Chapter 21, PostgreSQL actually does privilege management in terms of "roles". In this chapter, we consistently use *database user* to mean "role with the LOGIN privilege".

Authentication is the process by which the database server establishes the identity of the client, and by extension determines whether the client application (or the user who runs the client application) is permitted to connect with the database user name that was requested.

PostgreSQL offers a number of different client authentication methods. The method used to authenticate a particular client connection can be selected on the basis of (client) host address, database, and user.

PostgreSQL database user names are logically separate from user names of the operating system in which the server runs. If all the users of a particular server also have accounts on the server’s machine, it makes sense to assign database user names that match their operating system user names. However, a server that accepts remote connections might have many database users who have no local operating system account, and in such cases there need be no connection between database user names and OS user names.

## 20.1. The `pg_hba.conf` File

Client authentication is controlled by a configuration file, which traditionally is named `pg_hba.conf` and is stored in the database cluster’s data directory. (HBA stands for host-based authentication.) A default `pg_hba.conf` file is installed when the data directory is initialized by `initdb`. It is possible to place the authentication configuration file elsewhere, however; see the `hba_file` configuration parameter.

The general format of the `pg_hba.conf` file is a set of records, one per line. Blank lines are ignored, as is any text after the # comment character. Records cannot be continued across lines. A record is made up of a number of fields which are separated by spaces and/or tabs. Fields can contain white space if the field value is double-quoted. Quoting one of the keywords in a database, user, or address field (e.g., all or replication) makes the word lose its special meaning, and just match a database, user, or host with that name.

Each record specifies a connection type, a client IP address range (if relevant for the connection type), a database name, a user name, and the authentication method to be used for connections matching these parameters. The first record with a matching connection type, client address, requested database, and user name is used to perform authentication. There is no “fall-through” or “backup”; if one record is chosen and the authentication fails, subsequent records are not considered. If no record matches, access is denied.

A record can have one of the seven formats:

```
local  database user  auth-method  [auth-options]
host   database user  address   auth-method  [auth-options]
```
Chapter 20. Client Authentication

hostssl  database  user  address  auth-method  [auth-options]
hostnossl database  user  address  auth-method  [auth-options]
host       database  user  IP-address  IP-mask  auth-method  [auth-options]
hostssl   database  user  IP-address  IP-mask  auth-method  [auth-options]
hostnossl database  user  IP-address  IP-mask  auth-method  [auth-options]

The meaning of the fields is as follows:

local
This record matches connection attempts using Unix-domain sockets. Without a record of this
type, Unix-domain socket connections are disallowed.

host
This record matches connection attempts made using TCP/IP. host records match either SSL or
non-SSL connection attempts.

Note: Remote TCP/IP connections will not be possible unless the server is started with an
appropriate value for the listen_addresses configuration parameter, since the default behav-
ior is to listen for TCP/IP connections only on the local loopback address localhost.

hostssl
This record matches connection attempts made using TCP/IP, but only when the connection is
made with SSL encryption.

To make use of this option the server must be built with SSL support. Furthermore, SSL must be
enabled at server start time by setting the ssl configuration parameter (see Section 18.9 for more
information).

hostnossl
This record type has the opposite behavior of hostssl; it only matches connection attempts
made over TCP/IP that do not use SSL.

database
Specifies which database name(s) this record matches. The value all specifies that it matches
all databases. The value sameuser specifies that the record matches if the requested database
has the same name as the requested user. The value samerole specifies that the requested user
must be a member of the role with the same name as the requested database. (samegroup is an
obsolete but still accepted spelling of samerole.) Superusers are not considered to be members
of a role for the purposes of samerole unless they are explicitly members of the role, directly
or indirectly, and not just by virtue of being a superuser. The value replication specifies that
the record matches if a replication connection is requested (note that replication connections
do not specify any particular database). Otherwise, this is the name of a specific PostgreSQL
database. Multiple database names can be supplied by separating them with commas. A separate
file containing database names can be specified by preceding the file name with @.

user
Specifies which database user name(s) this record matches. The value all specifies that it
matches all users. Otherwise, this is either the name of a specific database user, or a group
name preceded by +. (Recall that there is no real distinction between users and groups in Post-
greSQL; a + mark really means “match any of the roles that are directly or indirectly members
of this role”, while a name without a + mark matches only that specific role.) For this purpose, a
superuser is only considered to be a member of a role if they are explicitly a member of the role, directly or indirectly, and not just by virtue of being a superuser. Multiple user names can be supplied by separating them with commas. A separate file containing user names can be specified by preceding the file name with \@.

Specifies the client machine address(es) that this record matches. This field can contain either a host name, an IP address range, or one of the special key words mentioned below.

An IP address range is specified using standard numeric notation for the range’s starting address, then a slash (/) and a CIDR mask length. The mask length indicates the number of high-order bits of the client IP address that must match. Bits to the right of this should be zero in the given IP address. There must not be any white space between the IP address, the /, and the CIDR mask length.

Typical examples of an IPv4 address range specified this way are 172.20.143.89/32 for a single host, or 172.20.143.0/24 for a small network, or 10.6.0.0/16 for a larger one. An IPv6 address range might look like ::1/128 for a single host (in this case the IPv6 loopback address) or fe80::7a31:c1ff:0000:0000/96 for a small network. 0.0.0.0/0 represents all IPv4 addresses, and ::/0 represents all IPv6 addresses. To specify a single host, use a mask length of 32 for IPv4 or 128 for IPv6. In a network address, do not omit trailing zeroes.

An entry given in IPv4 format will match only IPv4 connections, and an entry given in IPv6 format will match only IPv6 connections, even if the represented address is in the IPv4-in-IPv6 range. Note that entries in IPv6 format will be rejected if the system’s C library does not have support for IPv6 addresses.

You can also write all to match any IP address, samehost to match any of the server’s own IP addresses, or samenet to match any address in any subnet that the server is directly connected to.

If a host name is specified (anything that is not an IP address range or a special key word is treated as a host name), that name is compared with the result of a reverse name resolution of the client’s IP address (e.g., reverse DNS lookup, if DNS is used). Host name comparisons are case insensitive. If there is a match, then a forward name resolution (e.g., forward DNS lookup) is performed on the host name to check whether any of the addresses it resolves to are equal to the client’s IP address. If both directions match, then the entry is considered to match. (The host name that is used in pg_hba.conf should be the one that address-to-name resolution of the client’s IP address returns, otherwise the line won’t be matched. Some host name databases allow associating an IP address with multiple host names, but the operating system will only return one host name when asked to resolve an IP address.)

A host name specification that starts with a dot (.) matches a suffix of the actual host name. So .example.com would match foo.example.com (but not just example.com).

When host names are specified in pg_hba.conf, you should make sure that name resolution is reasonably fast. It can be of advantage to set up a local name resolution cache such as nscd. Also, you may wish to enable the configuration parameter log_hostname to see the client’s host name instead of the IP address in the log.

This field only applies to host, hostssl, and hostnossl records.
Chapter 20. Client Authentication

Users sometimes wonder why host names are handled in this seemingly complicated way, with two name resolutions including a reverse lookup of the client’s IP address. This complicates use of the feature in case the client’s reverse DNS entry is not set up or yields some undesirable host name. It is done primarily for efficiency: this way, a connection attempt requires at most two resolver lookups, one reverse and one forward. If there is a resolver problem with some address, it becomes only that client’s problem. A hypothetical alternative implementation that only did forward lookups would have to resolve every host name mentioned in `pg_hba.conf` during every connection attempt. That could be quite slow if many names are listed. And if there is a resolver problem with one of the host names, it becomes everyone’s problem.

Also, a reverse lookup is necessary to implement the suffix matching feature, because the actual client host name needs to be known in order to match it against the pattern.

Note that this behavior is consistent with other popular implementations of host name-based access control, such as the Apache HTTP Server and TCP Wrappers.

**IP-address**

**IP-mask**

These two fields can be used as an alternative to the `IP-address/mask-length` notation. Instead of specifying the mask length, the actual mask is specified in a separate column. For example, `255.0.0.0` represents an IPv4 CIDR mask length of 8, and `255.255.255.255` represents a CIDR mask length of 32.

These fields only apply to `host`, `hostssl`, and `hostnossl` records.

**auth-method**

Specifies the authentication method to use when a connection matches this record. The possible choices are summarized here; details are in Section 20.3.

**trust**

Allow the connection unconditionally. This method allows anyone that can connect to the PostgreSQL database server to login as any PostgreSQL user they wish, without the need for a password or any other authentication. See Section 20.3.1 for details.

**reject**

Reject the connection unconditionally. This is useful for “filtering out” certain hosts from a group, for example a `reject` line could block a specific host from connecting, while a later line allows the remaining hosts in a specific network to connect.

**md5**

Require the client to supply a double-MD5-hashed password for authentication. See Section 20.3.2 for details.

**password**

Require the client to supply an unencrypted password for authentication. Since the password is sent in clear text over the network, this should not be used on untrusted networks. See Section 20.3.2 for details.
Chapter 20. Client Authentication

Use GSSAPI to authenticate the user. This is only available for TCP/IP connections. See Section 20.3.3 for details.

Use SSPI to authenticate the user. This is only available on Windows. See Section 20.3.4 for details.

Obtain the operating system user name of the client by contacting the ident server on the client and check if it matches the requested database user name. Ident authentication can only be used on TCP/IP connections. When specified for local connections, peer authentication will be used instead. See Section 20.3.5 for details.

Obtain the client’s operating system user name from the operating system and check if it matches the requested database user name. This is only available for local connections. See Section 20.3.6 for details.

Authenticate using an LDAP server. See Section 20.3.7 for details.

Authenticate using a RADIUS server. See Section 20.3.8 for details.

Authenticate using SSL client certificates. See Section 20.3.9 for details.

Authenticate using the Pluggable Authentication Modules (PAM) service provided by the operating system. See Section 20.3.10 for details.

Authenticate using the BSD Authentication service provided by the operating system. See Section 20.3.11 for details.

auth-options

After the auth-method field, there can be field(s) of the form name=value that specify options for the authentication method. Details about which options are available for which authentication methods appear below.

In addition to the method-specific options listed below, there is one method-independent authentication option clientcert, which can be specified in any hostssl record. When set to 1, this option requires the client to present a valid (trusted) SSL certificate, in addition to the other requirements of the authentication method.

Files included by @ constructs are read as lists of names, which can be separated by either whitespace or commas. Comments are introduced by #, just as in pg_hba.conf, and nested @ constructs are allowed. Unless the file name following @ is an absolute path, it is taken to be relative to the directory containing the referencing file.
Since the `pg_hba.conf` records are examined sequentially for each connection attempt, the order of the records is significant. Typically, earlier records will have tight connection match parameters and weaker authentication methods, while later records will have looser match parameters and stronger authentication methods. For example, one might wish to use `trust` authentication for local TCP/IP connections but require a password for remote TCP/IP connections. In this case a record specifying `trust` authentication for connections from 127.0.0.1 would appear before a record specifying password authentication for a wider range of allowed client IP addresses.

The `pg_hba.conf` file is read on start-up and when the main server process receives a SIGHUP signal. If you edit the file on an active system, you will need to signal the postmaster (using `pg_ctl` `reload`, calling the SQL function `pg_reload_conf()`, or using `kill -HUP`) to make it re-read the file.

**Tip:** To connect to a particular database, a user must not only pass the `pg_hba.conf` checks, but must have the `CONNECT` privilege for the database. If you wish to restrict which users can connect to which databases, it's usually easier to control this by granting/revoking `CONNECT` privilege than to put the rules in `pg_hba.conf` entries.

Some examples of `pg_hba.conf` entries are shown in Example 20-1. See the next section for details on the different authentication methods.

**Example 20-1. Example pg_hba.conf Entries**

```
# Allow any user on the local system to connect to any database with
# any database user name using Unix-domain sockets (the default for local
# connections).
#
# TYPE DATABASE USER ADDRESS METHOD
local all all trust

# The same using local loopback TCP/IP connections.
#
# TYPE DATABASE USER ADDRESS METHOD
host all all 127.0.0.1/32 trust

# The same as the previous line, but using a separate netmask column
#
# TYPE DATABASE USER IP-ADDRESS IP-MASK METHOD
host all all 127.0.0.1 255.255.255.255 trust

# The same over IPv6.
#
# TYPE DATABASE USER ADDRESS METHOD
host all all ::1/128 trust

# The same using a host name (would typically cover both IPv4 and IPv6).
#
# TYPE DATABASE USER ADDRESS METHOD
host all all localhost trust

# Allow any user from any host with IP address 192.168.93.x to connect
# to database "postgres" as the same user name that ident reports for
# the connection (typically the operating system user name).
#
# TYPE DATABASE USER ADDRESS METHOD
```
Chapter 20. Client Authentication

host postgres all 192.168.93.0/24 ident

# Allow any user from host 192.168.12.10 to connect to database
# "postgres" if the user's password is correctly supplied.
#
# TYPE DATABASE USER ADDRESS METHOD
host postgres all 192.168.12.10/32 md5

# Allow any user from hosts in the example.com domain to connect to
database if the user's password is correctly supplied.
#
# TYPE DATABASE USER ADDRESS METHOD
host all all .example.com md5

# In the absence of preceding "host" lines, these two lines will
# reject all connections from 192.168.54.1 (since that entry will be
# matched first), but allow GSSAPI connections from anywhere else
# on the Internet. The zero mask causes no bits of the host IP
# address to be considered, so it matches any host.
#
# TYPE DATABASE USER ADDRESS METHOD
host all all 192.168.54.1/32 reject
host all all 0.0.0.0/0 gss

# Allow users from 192.168.x.x hosts to connect to any database, if
# they pass the ident check. If, for example, ident says the user is
# "bryanh" and he requests to connect as PostgreSQL user "guest1", the
# connection is allowed if there is an entry in pg_ident.conf for map
# "omicron" that says "bryanh" is allowed to connect as "guest1".
#
# TYPE DATABASE USER ADDRESS METHOD
host all all 192.168.0.0/16 ident map=omicron

# If these are the only three lines for local connections, they will
# allow local users to connect only to their own databases (databases
# with the same name as their database user name) except for administrators
# and members of role "support", who can connect to all databases. The file
# $PGDATA/admins contains a list of names of administrators. Passwords
# are required in all cases.
#
# TYPE DATABASE USER ADDRESS METHOD
local sameuser all md5
local all @admins md5
local all +support md5

# The last two lines above can be combined into a single line:
local all @admins,+support md5

# The database column can also use lists and file names:
local db1,db2,@demodbs all md5
20.2. User Name Maps

When using an external authentication system such as Ident or GSSAPI, the name of the operating system user that initiated the connection might not be the same as the database user (role) that is to be used. In this case, a user name map can be applied to map the operating system user name to a database user. To use user name mapping, specify `map=map-name` in the options field in `pg_hba.conf`. This option is supported for all authentication methods that receive external user names. Since different mappings might be needed for different connections, the name of the map to be used is specified in the `map-name` parameter in `pg_hba.conf` to indicate which map to use for each individual connection.

User name maps are defined in the ident map file, which by default is named `pg_ident.conf` and is stored in the cluster's data directory. (It is possible to place the map file elsewhere, however; see the `ident_file` configuration parameter.) The ident map file contains lines of the general form:

```
map-name  system-username  database-username
```

Comments and whitespace are handled in the same way as in `pg_hba.conf`. The `map-name` is an arbitrary name that will be used to refer to this mapping in `pg_hba.conf`. The other two fields specify an operating system user name and a matching database user name. The same `map-name` can be used repeatedly to specify multiple user-mappings within a single map.

There is no restriction regarding how many database users a given operating system user can correspond to, nor vice versa. Thus, entries in a map should be thought of as meaning "this operating system user is allowed to connect as this database user", rather than implying that they are equivalent. The connection will be allowed if there is any map entry that pairs the user name obtained from the external authentication system with the database user name that the user has requested to connect as.

If the `system-username` field starts with a slash (/), the remainder of the field is treated as a regular expression. (See Section 9.7.3.1 for details of PostgreSQL’s regular expression syntax.) The regular expression can include a single capture, or parenthesized subexpression, which can then be referenced in the `database-username` field as \1 (backslash-one). This allows the mapping of multiple user names in a single line, which is particularly useful for simple syntax substitutions. For example, these entries

```
mymap /^.\*@mydomain\.com$\ \1
mymap /^.\*@otherdomain\.com$\ guest
```

will remove the domain part for users with system user names that end with `@mydomain.com`, and allow any user whose system name ends with `@otherdomain.com` to log in as `guest`.

**Tip:** Keep in mind that by default, a regular expression can match just part of a string. It’s usually wise to use `^` and `$`, as shown in the above example, to force the match to be to the entire system user name.

The `pg_ident.conf` file is read on start-up and when the main server process receives a SIGHUP signal. If you edit the file on an active system, you will need to signal the postmaster (using `pg_ctl reload`, calling the SQL function `pg_reload_conf()`, or using `kill -HUP`) to make it re-read the file.

A `pg_ident.conf` file that could be used in conjunction with the `pg_hba.conf` file in Example 20-1 is shown in Example 20-2. In this example, anyone logged in to a machine on the 192.168 network that does not have the operating system user name `bryan`, `ann`, or `robert` would not be granted access. Unix user `robert` would only be allowed access when he tries to connect as PostgreSQL user 582.
bob, not as robert or anyone else. ann would only be allowed to connect as ann. User bryanh would be allowed to connect as either bryanh or as guest1.

Example 20-2. An Example pg_ident.conf File

# MAPNAME SYSTEM-USERNAME PG-USERNAME
omicron bryanh bryanh
omicron ann ann
# bob has user name robert on these machines
omicron robert bob
# bryanh can also connect as guest1
omicron bryanh guest1

20.3. Authentication Methods

The following subsections describe the authentication methods in more detail.

20.3.1. Trust Authentication

When trust authentication is specified, PostgreSQL assumes that anyone who can connect to the server is authorized to access the database with whatever database user name they specify (even superuser names). Of course, restrictions made in the database and user columns still apply. This method should only be used when there is adequate operating-system-level protection on connections to the server.

trust authentication is appropriate and very convenient for local connections on a single-user workstation. It is usually not appropriate by itself on a multiuser machine. However, you might be able to use trust even on a multiuser machine, if you restrict access to the server’s Unix-domain socket file using file-system permissions. To do this, set the unix_socket_permissions (and possibly unix_socket_group) configuration parameters as described in Section 19.3. Or you could set the unix_socket_directories configuration parameter to place the socket file in a suitably restricted directory.

Setting file-system permissions only helps for Unix-socket connections. Local TCP/IP connections are not restricted by file-system permissions. Therefore, if you want to use file-system permissions for local security, remove the host ... 127.0.0.1 ... line from pg_hba.conf, or change it to a non-trust authentication method.

trust authentication is only suitable for TCP/IP connections if you trust every user on every machine that is allowed to connect to the server by the pg_hba.conf lines that specify trust. It is seldom reasonable to use trust for any TCP/IP connections other than those from localhost (127.0.0.1).

20.3.2. Password Authentication

The password-based authentication methods are md5 and password. These methods operate similarly except for the way that the password is sent across the connection, namely MD5-hashed and clear-text respectively.

If you are at all concerned about password “sniffing” attacks then md5 is preferred. Plain password should always be avoided if possible. However, md5 cannot be used with the db_user_namespace
Chapter 20. Client Authentication

feature. If the connection is protected by SSL encryption then password can be used safely (though SSL certificate authentication might be a better choice if one is depending on using SSL).

PostgreSQL database passwords are separate from operating system user passwords. The password for each database user is stored in the pg_authid system catalog. Passwords can be managed with the SQL commands CREATE USER and ALTER ROLE, e.g., CREATE USER foo WITH PASSWORD 'secret'. If no password has been set up for a user, the stored password is null and password authentication will always fail for that user.

20.3.3. GSSAPI Authentication

GSSAPI is an industry-standard protocol for secure authentication defined in RFC 2743. PostgreSQL supports GSSAPI with Kerberos authentication according to RFC 1964. GSSAPI provides automatic authentication (single sign-on) for systems that support it. The authentication itself is secure, but the data sent over the database connection will be sent unencrypted unless SSL is used.

GSSAPI support has to be enabled when PostgreSQL is built; see Chapter 16 for more information.

When GSSAPI uses Kerberos, it uses a standard principal in the format servicename/hostname@realm. The PostgreSQL server will accept any principal that is included in the keytab used by the server, but care needs to be taken to specify the correct principal details when making the connection from the client using the krbsrvname connection parameter. (See also Section 32.1.2.) The installation default can be changed from the default postgres at build time using ./configure --with-krb-srvname=whatever. In most environments, this parameter never needs to be changed. Some Kerberos implementations might require a different service name, such as Microsoft Active Directory which requires the service name to be in upper case (POSTGRES).

hostname is the fully qualified host name of the server machine. The service principal’s realm is the preferred realm of the server machine.

Client principals can be mapped to different PostgreSQL database user names with pg_ident.conf. For example, pgusername@realm could be mapped to just pgusername. Alternatively, you can use the full username@realm principal as the role name in PostgreSQL without any mapping.

PostgreSQL also supports a parameter to strip the realm from the principal. This method is supported for backwards compatibility and is strongly discouraged as it is then impossible to distinguish different users with the same user name but coming from different realms. To enable this, set include_realm to 0. For simple single-realm installations, doing that combined with setting the krb_realm parameter (which checks that the principal’s realm matches exactly what is in the krb_realm parameter) is still secure; but this is a less capable approach compared to specifying an explicit mapping in pg_ident.conf.

Make sure that your server keytab file is readable (and preferably only readable, not writable) by the PostgreSQL server account. (See also Section 18.1.) The location of the key file is specified by the krb_server_keyfile configuration parameter. The default is /usr/local/pgsql/etc/krb5.keytab (or whatever directory was specified as sysconfdir at build time). For security reasons, it is recommended to use a separate keytab just for the PostgreSQL server rather than opening up permissions on the system keytab file.

The keytab file is generated by the Kerberos software; see the Kerberos documentation for details. The following example is for MIT-compatible Kerberos 5 implementations:

```
admin% kadmin -randkey postgres/server.my.domain.org
admin% ktadd -k krb5.keytab postgres/server.my.domain.org
```
When connecting to the database make sure you have a ticket for a principal matching the requested database user name. For example, for database user name `fred`, principal `fred@EXAMPLE.COM` would be able to connect. To also allow principal `fred/users.example.com@EXAMPLE.COM`, use a user name map, as described in Section 20.2.

The following configuration options are supported for GSSAPI:

`include_realm`

If set to 0, the realm name from the authenticated user principal is stripped off before being passed through the user name mapping (Section 20.2). This is discouraged and is primarily available for backwards compatibility, as it is not secure in multi-realm environments unless `krb_realm` is also used. It is recommended to leave `include_realm` set to the default (1) and to provide an explicit mapping in `pg_ident.conf` to convert principal names to PostgreSQL user names.

`map`

Allows for mapping between system and database user names. See Section 20.2 for details. For a GSSAPI/Kerberos principal, such as `username@EXAMPLE.COM` (or, less commonly, `username/hostbased@EXAMPLE.COM`), the user name used for mapping is `username@EXAMPLE.COM` (or `username/hostbased@EXAMPLE.COM`, respectively), unless `include_realm` has been set to 0, in which case `username` (or `username/hostbased`) is what is seen as the system user name when mapping.

`krb_realm`

Sets the realm to match user principal names against. If this parameter is set, only users of that realm will be accepted. If it is not set, users of any realm can connect, subject to whatever user name mapping is done.

### 20.3.4. SSPI Authentication

SSPI is a Windows technology for secure authentication with single sign-on. PostgreSQL will use SSPI in `negotiate` mode, which will use Kerberos when possible and automatically fall back to NTLM in other cases. SSPI authentication only works when both server and client are running Windows, or, on non-Windows platforms, when GSSAPI is available.

When using Kerberos authentication, SSPI works the same way GSSAPI does; see Section 20.3.3 for details.

The following configuration options are supported for SSPI:

`include_realm`

If set to 0, the realm name from the authenticated user principal is stripped off before being passed through the user name mapping (Section 20.2). This is discouraged and is primarily available for backwards compatibility, as it is not secure in multi-realm environments unless `krb_realm` is also used. It is recommended to leave `include_realm` set to the default (1) and to provide an explicit mapping in `pg_ident.conf` to convert principal names to PostgreSQL user names.

`compat_realm`

If set to 1, the domain’s SAM-compatible name (also known as the NetBIOS name) is used for the `include_realm` option. This is the default. If set to 0, the true realm name from the
Kerberos user principal name is used.

Do not disable this option unless your server runs under a domain account (this includes virtual service accounts on a domain member system) and all clients authenticating through SSPI are also using domain accounts, or authentication will fail.

`upn_username`

If this option is enabled along with `compat_realm`, the user name from the Kerberos UPN is used for authentication. If it is disabled (the default), the SAM-compatible user name is used. By default, these two names are identical for new user accounts.

Note that libpq uses the SAM-compatible name if no explicit user name is specified. If you use libpq or a driver based on it, you should leave this option disabled or explicitly specify user name in the connection string.

`map`

Allows for mapping between system and database user names. See Section 20.2 for details. For a SSPI/Kerberos principal, such as `username@EXAMPLE.COM` (or, less commonly, `username/hostbased@EXAMPLE.COM`), the user name used for mapping is `username@EXAMPLE.COM` (or `username/hostbased@EXAMPLE.COM`, respectively), unless `include_realm` has been set to 0, in which case `username` (or `username/hostbased`) is what is seen as the system user name when mapping.

`krb.realm`

Sets the realm to match user principal names against. If this parameter is set, only users of that realm will be accepted. If it is not set, users of any realm can connect, subject to whatever user name mapping is done.

### 20.3.5. Ident Authentication

The ident authentication method works by obtaining the client’s operating system user name from an ident server and using it as the allowed database user name (with an optional user name mapping). This is only supported on TCP/IP connections.

**Note:** When ident is specified for a local (non-TCP/IP) connection, peer authentication (see Section 20.3.6) will be used instead.

The following configuration options are supported for ident:

`map`

Allows for mapping between system and database user names. See Section 20.2 for details.

The “Identification Protocol” is described in RFC 1413. Virtually every Unix-like operating system ships with an ident server that listens on TCP port 113 by default. The basic functionality of an ident server is to answer questions like “What user initiated the connection that goes out of your port \(x\) and connects to my port \(y\)?” Since PostgreSQL knows both \(x\) and \(y\) when a physical connection is established, it can interrogate the ident server on the host of the connecting client and can theoretically determine the operating system user for any given connection.
Chapter 20. Client Authentication

The drawback of this procedure is that it depends on the integrity of the client: if the client machine is untrusted or compromised, an attacker could run just about any program on port 113 and return any user name they choose. This authentication method is therefore only appropriate for closed networks where each client machine is under tight control and where the database and system administrators operate in close contact. In other words, you must trust the machine running the ident server. Heed the warning:

The Identification Protocol is not intended as an authorization or access control protocol.

—RFC 1413

Some ident servers have a nonstandard option that causes the returned user name to be encrypted, using a key that only the originating machine’s administrator knows. This option must not be used when using the ident server with PostgreSQL, since PostgreSQL does not have any way to decrypt the returned string to determine the actual user name.

20.3.6. Peer Authentication

The peer authentication method works by obtaining the client’s operating system user name from the kernel and using it as the allowed database user name (with optional user name mapping). This method is only supported on local connections.

The following configuration options are supported for peer:

map

Allows for mapping between system and database user names. See Section 20.2 for details.

Peer authentication is only available on operating systems providing the getpeereid() function, the SO_PEERCRED socket parameter, or similar mechanisms. Currently that includes Linux, most flavors of BSD including OS X, and Solaris.

20.3.7. LDAP Authentication

This authentication method operates similarly to password except that it uses LDAP as the password verification method. LDAP is used only to validate the user name/password pairs. Therefore the user must already exist in the database before LDAP can be used for authentication.

LDAP authentication can operate in two modes. In the first mode, which we will call the simple bind mode, the server will bind to the distinguished name constructed as prefix username suffix. Typically, the prefix parameter is used to specify cn=, or DOMAIN\ in an Active Directory environment. suffix is used to specify the remaining part of the DN in a non-Active Directory environment.

In the second mode, which we will call the search+bind mode, the server first binds to the LDAP directory with a fixed user name and password, specified with ldapbinddn and ldapbindpasswd, and performs a search for the user trying to log in to the database. If no user and password is configured, an anonymous bind will be attempted to the directory. The search will be performed over the subtree at ldappasedn, and will try to do an exact match of the attribute specified in ldapsearchattribute. Once the user has been found in this search, the server disconnects and re-binds to the directory as this user, using the password specified by the client, to verify that the login is correct. This mode is the same as that used by LDAP authentication schemes in other software, such as Apache mod_authnz_ldap and pam_ldap. This method allows for significantly more flexibility in where the user objects are located in the directory, but will cause two separate connections to the LDAP server to be made.
Chapter 20. Client Authentication

The following configuration options are used in both modes:

**ldapserver**

Names or IP addresses of LDAP servers to connect to. Multiple servers may be specified, separated by spaces.

**ldapport**

Port number on LDAP server to connect to. If no port is specified, the LDAP library’s default port setting will be used.

**ldaptls**

Set to 1 to make the connection between PostgreSQL and the LDAP server use TLS encryption. Note that this only encrypts the traffic to the LDAP server — the connection to the client will still be unencrypted unless SSL is used.

The following options are used in simple bind mode only:

**ldapprefix**

String to prepend to the user name when forming the DN to bind as, when doing simple bind authentication.

**ldapsuffix**

String to append to the user name when forming the DN to bind as, when doing simple bind authentication.

The following options are used in search+bind mode only:

**ldapbasedn**

Root DN to begin the search for the user in, when doing search+bind authentication.

**ldapbinddn**

DN of user to bind to the directory with to perform the search when doing search+bind authentication.

**ldapbindpasswd**

Password for user to bind to the directory with to perform the search when doing search+bind authentication.

**ldapsearchattribute**

Attribute to match against the user name in the search when doing search+bind authentication. If no attribute is specified, the `uid` attribute will be used.

**ldapurl**

An RFC 4516 LDAP URL. This is an alternative way to write some of the other LDAP options in a more compact and standard form. The format is

```
ldap://host[:port]/basedn[?attribute][?[scope]]
```

*scope* must be one of `base`, `one`, `sub`, typically the latter. Only one attribute is used, and some other components of standard LDAP URLs such as filters and extensions are not supported.

For non-anonymous binds, `ldapbinddn` and `ldapbindpasswd` must be specified as separate options.

To use encrypted LDAP connections, the `ldaptls` option has to be used in addition to `ldapurl`. The `ldaps` URL scheme (direct SSL connection) is not supported.

LDAP URLs are currently only supported with OpenLDAP, not on Windows.
It is an error to mix configuration options for simple bind with options for search+bind.

Here is an example for a simple-bind LDAP configuration:

```plaintext
host ... ldap ldapserver=ldap.example.net ldapprefix="cn=" ldapsuffix="", dc=example, dc=
```

When a connection to the database server as database user `someuser` is requested, PostgreSQL will attempt to bind to the LDAP server using the DN `cn=someuser, dc=example, dc=net` and the password provided by the client. If that connection succeeds, the database access is granted.

Here is an example for a search+bind configuration:

```plaintext
host ... ldap ldapserver=ldap.example.net ldapbasedn="dc=example, dc=net" ldapsearchattribute=uid
```

When a connection to the database server as database user `someuser` is requested, PostgreSQL will attempt to bind anonymously (since `ldapbinddn` was not specified) to the LDAP server, perform a search for `(uid=someuser)` under the specified base DN. If an entry is found, it will then attempt to bind using that found information and the password supplied by the client. If that second connection succeeds, the database access is granted.

Here is the same search+bind configuration written as a URL:

```plaintext
host ... ldap ldapurl="ldap://ldap.example.net/dc=example,dc=net?uid?sub"
```

Some other software that supports authentication against LDAP uses the same URL format, so it will be easier to share the configuration.

**Tip:** Since LDAP often uses commas and spaces to separate the different parts of a DN, it is often necessary to use double-quoted parameter values when configuring LDAP options, as shown in the examples.

### 20.3.8. RADIUS Authentication

This authentication method operates similarly to password except that it uses RADIUS as the password verification method. RADIUS is used only to validate the user name/password pairs. Therefore the user must already exist in the database before RADIUS can be used for authentication.

When using RADIUS authentication, an Access Request message will be sent to the configured RADIUS server. This request will be of type `Authenticate Only`, and include parameters for `user name`, `password` (encrypted) and `NAS Identifier`. The request will be encrypted using a secret shared with the server. The RADIUS server will respond to this server with either `Access Accept` or `Access Reject`. There is no support for RADIUS accounting.

The following configuration options are supported for RADIUS:

- `radiusserver`

  The name or IP address of the RADIUS server to connect to. This parameter is required.
radiussecret
The shared secret used when talking securely to the RADIUS server. This must have exactly the
same value on the PostgreSQL and RADIUS servers. It is recommended that this be a string of
at least 16 characters. This parameter is required.

Note: The encryption vector used will only be cryptographically strong if PostgreSQL is built
with support for OpenSSL. In other cases, the transmission to the RADIUS server should
only be considered obfuscated, not secured, and external security measures should be ap-
plied if necessary.

radiusport
The port number on the RADIUS server to connect to. If no port is specified, the default port
1812 will be used.

radiusidentifier
The string used as NAS Identifier in the RADIUS requests. This parameter can be used
as a second parameter identifying for example which database user the user is attempting to
authenticate as, which can be used for policy matching on the RADIUS server. If no identifier is
specified, the default postgresql will be used.

20.3.9. Certificate Authentication
This authentication method uses SSL client certificates to perform authentication. It is therefore only
available for SSL connections. When using this authentication method, the server will require that
the client provide a valid, trusted certificate. No password prompt will be sent to the client. The cn
(Common Name) attribute of the certificate will be compared to the requested database user name,
and if they match the login will be allowed. User name mapping can be used to allow cn to be different
from the database user name.
The following configuration options are supported for SSL certificate authentication:

map
Allows for mapping between system and database user names. See Section 20.2 for details.

In a pg_hba.conf record specifying certificate authentication, the authentication option
clientcert is assumed to be 1, and it cannot be turned off since a client certificate is necessary for
this method. What the cert method adds to the basic clientcert certificate validity test is a check
that the cn attribute matches the database user name.

20.3.10. PAM Authentication
This authentication method operates similarly to password except that it uses PAM (Pluggable
Authentication Modules) as the authentication mechanism. The default PAM service name is
postgresql. PAM is used only to validate user name/password pairs and optionally the connected
remote host name or IP address. Therefore the user must already exist in the database before PAM can be used for authentication. For more information about PAM, please read the Linux-PAM Page.

The following configuration options are supported for PAM:

```plaintext
pamservice

PAM service name.

pam_use_hostname

Determines whether the remote IP address or the host name is provided to PAM modules through the PAM_RHOST item. By default, the IP address is used. Set this option to 1 to use the resolved host name instead. Host name resolution can lead to login delays. (Most PAM configurations don’t use this information, so it is only necessary to consider this setting if a PAM configuration was specifically created to make use of it.)

Note: If PAM is set up to read /etc/shadow, authentication will fail because the PostgreSQL server is started by a non-root user. However, this is not an issue when PAM is configured to use LDAP or other authentication methods.

20.3.11. BSD Authentication

This authentication method operates similarly to password except that it uses BSD Authentication to verify the password. BSD Authentication is used only to validate user name/password pairs. Therefore the user’s role must already exist in the database before BSD Authentication can be used for authentication. The BSD Authentication framework is currently only available on OpenBSD.

BSD Authentication in PostgreSQL uses the auth-postgresql login type and authenticates with the postgresql login class if that’s defined in login.conf. By default that login class does not exist, and PostgreSQL will use the default login class.

Note: To use BSD Authentication, the PostgreSQL user account (that is, the operating system user running the server) must first be added to the auth group. The auth group exists by default on OpenBSD systems.

20.4. Authentication Problems

Authentication failures and related problems generally manifest themselves through error messages like the following:

```
FATAL: no pg_hba.conf entry for host "123.123.123.123", user "andym", database "testdb"
```

This is what you are most likely to get if you succeed in contacting the server, but it does not want to talk to you. As the message suggests, the server refused the connection request because it found no matching entry in its pg_hba.conf configuration file.

Chapter 20. Client Authentication

FATAL: password authentication failed for user "andym"

Messages like this indicate that you contacted the server, and it is willing to talk to you, but not until you pass the authorization method specified in the `pg_hba.conf` file. Check the password you are providing, or check your Kerberos or ident software if the complaint mentions one of those authentication types.

FATAL: user "andym" does not exist

The indicated database user name was not found.

FATAL: database "testdb" does not exist

The database you are trying to connect to does not exist. Note that if you do not specify a database name, it defaults to the database user name, which might or might not be the right thing.

**Tip:** The server log might contain more information about an authentication failure than is reported to the client. If you are confused about the reason for a failure, check the server log.
Chapter 21. Database Roles

PostgreSQL manages database access permissions using the concept of roles. A role can be thought of as either a database user, or a group of database users, depending on how the role is set up. Roles can own database objects (for example, tables and functions) and can assign privileges on those objects to other roles to control who has access to which objects. Furthermore, it is possible to grant membership in a role to another role, thus allowing the member role to use privileges assigned to another role.

The concept of roles subsumes the concepts of “users” and “groups”. In PostgreSQL versions before 8.1, users and groups were distinct kinds of entities, but now there are only roles. Any role can act as a user, a group, or both.

This chapter describes how to create and manage roles. More information about the effects of role privileges on various database objects can be found in Section 5.6.

21.1. Database Roles

Database roles are conceptually completely separate from operating system users. In practice it might be convenient to maintain a correspondence, but this is not required. Database roles are global across a database cluster installation (and not per individual database). To create a role use the CREATE ROLE SQL command:

```
CREATE ROLE name;
```

`name` follows the rules for SQL identifiers: either unadorned without special characters, or double-quoted. (In practice, you will usually want to add additional options, such as LOGIN, to the command. More details appear below.) To remove an existing role, use the analogous DROP ROLE command:

```
DROP ROLE name;
```

For convenience, the programs createuser and dropuser are provided as wrappers around these SQL commands that can be called from the shell command line:

```
createuser name
dropuser name
```

To determine the set of existing roles, examine the `pg_roles` system catalog, for example

```
SELECT rolname FROM pg_roles;
```

The `psql` program’s `\du` meta-command is also useful for listing the existing roles.

In order to bootstrap the database system, a freshly initialized system always contains one predefined role. This role is always a “superuser”, and by default (unless altered when running `initdb`) it will have the same name as the operating system user that initialized the database cluster. Customarily, this role will be named `postgres`. In order to create more roles you first have to connect as this initial role.

Every connection to the database server is made using the name of some particular role, and this role determines the initial access privileges for commands issued in that connection. The role name to use for a particular database connection is indicated by the client that is initiating the connection request in an application-specific fashion. For example, the `psql` program uses the `-U` command line option.
to indicate the role to connect as. Many applications assume the name of the current operating system user by default (including `createuser` and `psql`). Therefore it is often convenient to maintain a naming correspondence between roles and operating system users.

The set of database roles a given client connection can connect as is determined by the client authentication setup, as explained in Chapter 20. (Thus, a client is not limited to connect as the role matching its operating system user, just as a person’s login name need not match his or her real name.) Since the role identity determines the set of privileges available to a connected client, it is important to carefully configure privileges when setting up a multiuser environment.

### 21.2. Role Attributes

A database role can have a number of attributes that define its privileges and interact with the client authentication system.

**login privilege**

Only roles that have the `LOGIN` attribute can be used as the initial role name for a database connection. A role with the `LOGIN` attribute can be considered the same as a "database user". To create a role with login privilege, use either:

```
CREATE ROLE name LOGIN;
CREATE USER name;
```

*(CREATE USER is equivalent to CREATE ROLE except that CREATE USER assumes LOGIN by default, while CREATE ROLE does not.)*

**superuser status**

A database superuser bypasses all permission checks, except the right to log in. This is a dangerous privilege and should not be used carelessly; it is best to do most of your work as a role that is not a superuser. To create a new database superuser, use `CREATE ROLE name SUPERUSER`. You must do this as a role that is already a superuser.

**database creation**

A role must be explicitly given permission to create databases (except for superusers, since those bypass all permission checks). To create such a role, use `CREATE ROLE name CREATEDB`.

**role creation**

A role must be explicitly given permission to create more roles (except for superusers, since those bypass all permission checks). To create such a role, use `CREATE ROLE name CREATEROLE`. A role with `CREATEROLE` privilege can alter and drop other roles, too, as well as grant or revoke membership in them. However, to create, alter, drop, or change membership of a superuser role, superuser status is required; `CREATEROLE` is insufficient for that.

**initiating replication**

A role must explicitly be given permission to initiate streaming replication (except for superusers, since those bypass all permission checks). A role used for streaming replication must have `LOGIN` permission as well. To create such a role, use `CREATE ROLE name REPLICATION LOGIN`.

**password**

A password is only significant if the client authentication method requires the user to supply a password when connecting to the database. The `password` and `md5` authentication methods make use of passwords. Database passwords are separate from operating system passwords. Specify a password upon role creation with `CREATE ROLE name PASSWORD 'string'`. 

594
A role’s attributes can be modified after creation with `ALTER ROLE`. See the reference pages for the `CREATE ROLE` and `ALTER ROLE` commands for details.

**Tip:** It is good practice to create a role that has the `CREATEDB` and `CREATEROLE` privileges, but is not a superuser, and then use this role for all routine management of databases and roles. This approach avoids the dangers of operating as a superuser for tasks that do not really require it.

A role can also have role-specific defaults for many of the run-time configuration settings described in Chapter 19. For example, if for some reason you want to disable index scans (hint: not a good idea) anytime you connect, you can use:

```
ALTER ROLE myname SET enable_indexscan TO off;
```

This will save the setting (but not set it immediately). In subsequent connections by this role it will appear as though `SET enable_indexscan TO off` had been executed just before the session started. You can still alter this setting during the session; it will only be the default. To remove a role-specific default setting, use `ALTER ROLE rolename RESET varname`. Note that role-specific defaults attached to roles without `LOGIN` privilege are fairly useless, since they will never be invoked.

### 21.3. Role Membership

It is frequently convenient to group users together to ease management of privileges: that way, privileges can be granted to, or revoked from, a group as a whole. In PostgreSQL this is done by creating a role that represents the group, and then granting membership in the group role to individual user roles.

To set up a group role, first create the role:

```
CREATE ROLE name;
```

Typically a role being used as a group would not have the `LOGIN` attribute, though you can set it if you wish.

Once the group role exists, you can add and remove members using the `GRANT` and `REVOKE` commands:

```
GRANT group_role TO role1, ... ;
REVOKE group_role FROM role1, ... ;
```

You can grant membership to other group roles, too (since there isn’t really any distinction between group roles and non-group roles). The database will not let you set up circular membership loops. Also, it is not permitted to grant membership in a role to `PUBLIC`.

The members of a group role can use the privileges of the role in two ways. First, every member of a group can explicitly do `SET ROLE` to temporarily “become” the group role. In this state, the database session has access to the privileges of the group role rather than the original login role, and any database objects created are considered owned by the group role not the login role. Second, member roles that have the `INHERIT` attribute automatically have use of the privileges of roles of which they are members, including any privileges inherited by those roles. As an example, suppose we have done:

```
CREATE ROLE joe LOGIN INHERIT;
CREATE ROLE admin NOINHERIT;
```
CREATE ROLE wheel NOINHERIT;
GRANT admin TO joe;
GRANT wheel TO admin;

Immediately after connecting as role joe, a database session will have use of privileges granted directly to joe plus any privileges granted to admin, because joe “inherits” admin’s privileges. However, privileges granted to wheel are not available, because even though joe is indirectly a member of wheel, the membership is via admin which has the NOINHERIT attribute. After:

SET ROLE admin;

the session would have use of only those privileges granted to admin, and not those granted to joe. After:

SET ROLE wheel;

the session would have use of only those privileges granted to wheel, and not those granted to either joe or admin. The original privilege state can be restored with any of:

SET ROLE joe;
SET ROLE NONE;
RESET ROLE;

**Note:** The SET ROLE command always allows selecting any role that the original login role is directly or indirectly a member of. Thus, in the above example, it is not necessary to become admin before becoming wheel.

**Note:** In the SQL standard, there is a clear distinction between users and roles, and users do not automatically inherit privileges while roles do. This behavior can be obtained in PostgreSQL by giving roles being used as SQL roles the INHERIT attribute, while giving roles being used as SQL users the NOINHERIT attribute. However, PostgreSQL defaults to giving all roles the INHERIT attribute, for backward compatibility with pre-8.1 releases in which users always had use of permissions granted to groups they were members of.

The role attributes LOGIN, SUPERUSER, CREATEDB, and CREATEROLE can be thought of as special privileges, but they are never inherited as ordinary privileges on database objects are. You must actually SET ROLE to a specific role having one of these attributes in order to make use of the attribute. Continuing the above example, we might choose to grant CREATEDB and CREATEROLE to the admin role. Then a session connecting as role joe would not have these privileges immediately, only after doing SET ROLE admin.

To destroy a group role, use DROP ROLE:

DROP ROLE name;

Any memberships in the group role are automatically revoked (but the member roles are not otherwise affected).
21.4. Dropping Roles

Because roles can own database objects and can hold privileges to access other objects, dropping a role is often not just a matter of a quick DROP ROLE. Any objects owned by the role must first be dropped or reassigned to other owners; and any permissions granted to the role must be revoked.

Ownership of objects can be transferred one at a time using ALTER commands, for example:

```
ALTER TABLE bobs_table OWNER TO alice;
```

Alternatively, the REASSIGN OWNED command can be used to reassign ownership of all objects owned by the role-to-be-dropped to a single other role. Because REASSIGN OWNED cannot access objects in other databases, it is necessary to run it in each database that contains objects owned by the role. (Note that the first such REASSIGN OWNED will change the ownership of any shared-across-databases objects, that is databases or tablespaces, that are owned by the role-to-be-dropped.)

Once any valuable objects have been transferred to new owners, any remaining objects owned by the role-to-be-dropped can be dropped with the DROP OWNED command. Again, this command cannot access objects in other databases, so it is necessary to run it in each database that contains objects owned by the role. Also, DROP OWNED will not drop entire databases or tablespaces, so it is necessary to do that manually if the role owns any databases or tablespaces that have not been transferred to new owners.

DROP OWNED also takes care of removing any privileges granted to the target role for objects that do not belong to it. Because REASSIGN OWNED does not touch such objects, it’s typically necessary to run both REASSIGN OWNED and DROP OWNED (in that order!) to fully remove the dependencies of a role to be dropped.

In short then, the most general recipe for removing a role that has been used to own objects is:

```
REASSIGN OWNED BY doomed_role TO successor_role;
DROP OWNED BY doomed_role;
```

When not all owned objects are to be transferred to the same successor owner, it’s best to handle the exceptions manually and then perform the above steps to mop up.

If DROP ROLE is attempted while dependent objects still remain, it will issue messages identifying which objects need to be reassigned or dropped.

21.5. Default Roles

PostgreSQL provides a set of default roles which provide access to certain, commonly needed, privileged capabilities and information. Administrators can GRANT these roles to users and/or other roles in their environment, providing those users with access to the specified capabilities and information.

The default roles are described in Table 21-1. Note that the specific permissions for each of the default roles may change in the future as additional capabilities are added. Administrators should monitor the release notes for changes.

<table>
<thead>
<tr>
<th>Role</th>
<th>Allowed Access</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### Role Access

<table>
<thead>
<tr>
<th>Role</th>
<th>Allowed Access</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_signal_backend</td>
<td>Send signals to other backends (e.g.: cancel query, terminate).</td>
</tr>
</tbody>
</table>

Administrators can grant access to these roles to users using the GRANT command:

```sql
GRANT pg_signal_backend TO admin_user;
```

### 21.6. Function Security

Functions, triggers and row-level security policies allow users to insert code into the backend server that other users might execute unintentionally. Hence, these mechanisms permit users to “Trojan horse” others with relative ease. The strongest protection is tight control over who can define objects. Where that is infeasible, write queries referring only to objects having trusted owners. Remove from `search_path` the public schema and any other schemas that permit untrusted users to create objects.

Functions run inside the backend server process with the operating system permissions of the database server daemon. If the programming language used for the function allows unchecked memory accesses, it is possible to change the server’s internal data structures. Hence, among many other things, such functions can circumvent any system access controls. Function languages that allow such access are considered “untrusted”, and PostgreSQL allows only superusers to create functions written in those languages.
Chapter 22. Managing Databases

Every instance of a running PostgreSQL server manages one or more databases. Databases are therefore the topmost hierarchical level for organizing SQL objects (“database objects”). This chapter describes the properties of databases, and how to create, manage, and destroy them.

22.1. Overview

A database is a named collection of SQL objects (“database objects”). Generally, every database object (tables, functions, etc.) belongs to one and only one database. (However there are a few system catalogs, for example pg_database, that belong to a whole cluster and are accessible from each database within the cluster.) More accurately, a database is a collection of schemas and the schemas contain the tables, functions, etc. So the full hierarchy is: server, database, schema, table (or some other kind of object, such as a function).

When connecting to the database server, a client must specify in its connection request the name of the database it wants to connect to. It is not possible to access more than one database per connection. However, an application is not restricted in the number of connections it opens to the same or other databases. Databases are physically separated and access control is managed at the connection level. If one PostgreSQL server instance is to house projects or users that should be separate and for the most part unaware of each other, it is therefore recommended to put them into separate databases. If the projects or users are interrelated and should be able to use each other’s resources, they should be put in the same database but possibly into separate schemas. Schemas are a purely logical structure and who can access what is managed by the privilege system. More information about managing schemas is in Section 5.8.

Databases are created with the CREATE DATABASE command (see Section 22.2) and destroyed with the DROP DATABASE command (see Section 22.5). To determine the set of existing databases, examine the pg_database system catalog, for example

\[
\text{SELECT datname FROM pg_database;}
\]

The psql program’s \l meta-command and -l command-line option are also useful for listing the existing databases.

Note: The SQL standard calls databases “catalogs”, but there is no difference in practice.

22.2. Creating a Database

In order to create a database, the PostgreSQL server must be up and running (see Section 18.3).

Databases are created with the SQL command CREATE DATABASE:

\[
\text{CREATE DATABASE name;}
\]

where name follows the usual rules for SQL identifiers. The current role automatically becomes the owner of the new database. It is the privilege of the owner of a database to remove it later (which also removes all the objects in it, even if they have a different owner).

The creation of databases is a restricted operation. See Section 21.2 for how to grant permission.
Since you need to be connected to the database server in order to execute the `CREATE DATABASE` command, the question remains how the first database at any given site can be created. The first database is always created by the `initdb` command when the data storage area is initialized. (See Section 18.2.) This database is called `postgres`. So to create the first “ordinary” database you can connect to `postgres`.

A second database, `template1`, is also created during database cluster initialization. Whenever a new database is created within the cluster, `template1` is essentially cloned. This means that any changes you make in `template1` are propagated to all subsequently created databases. Because of this, avoid creating objects in `template1` unless you want them propagated to every newly created database. More details appear in Section 22.3.

As a convenience, there is a program you can execute from the shell to create new databases, `createdb`.

`createdb dbname` does no magic. It connects to the `postgres` database and issues the `CREATE DATABASE` command, exactly as described above. The createdb reference page contains the invocation details. Note that `createdb` without any arguments will create a database with the current user name.

**Note:** Chapter 20 contains information about how to restrict who can connect to a given database.

Sometimes you want to create a database for someone else, and have them become the owner of the new database, so they can configure and manage it themselves. To achieve that, use one of the following commands:

```
CREATE DATABASE dbname OWNER rolename;
```

from the SQL environment, or:

```
createdb -O rolename dbname
```

from the shell. Only the superuser is allowed to create a database for someone else (that is, for a role you are not a member of).

### 22.3. Template Databases

`CREATE DATABASE` actually works by copying an existing database. By default, it copies the standard system database named `template1`. Thus that database is the “template” from which new databases are made. If you add objects to `template1`, these objects will be copied into subsequently created user databases. This behavior allows site-local modifications to the standard set of objects in databases. For example, if you install the procedural language PL/Perl in `template1`, it will automatically be available in user databases without any extra action being taken when those databases are created.

There is a second standard system database named `template0`. This database contains the same data as the initial contents of `template1`, that is, only the standard objects predefined by your version of PostgreSQL. `template0` should never be changed after the database cluster has been initialized. By instructing `CREATE DATABASE` to copy `template0` instead of `template1`, you can create a “virgin” user database that contains none of the site-local additions in `template1`. This is particularly handy when restoring a `pg_dump` dump: the dump script should be restored in a virgin database to ensure
that one recreates the correct contents of the dumped database, without conflicting with objects that
might have been added to template1 later on.

Another common reason for copying template0 instead of template1 is that new encoding and
locale settings can be specified when copying template0, whereas a copy of template1 must use
the same settings it does. This is because template1 might contain encoding-specific or locale-
specific data, while template0 is known not to.

To create a database by copying template0, use:

```
CREATE DATABASE dbname TEMPLATE template0;
```

from the SQL environment, or:

```
createdb -T template0 dbname
```

from the shell.

It is possible to create additional template databases, and indeed one can copy any database in a
cluster by specifying its name as the template for CREATE DATABASE. It is important to understand,
however, that this is not (yet) intended as a general-purpose "COPY DATABASE" facility. The principal
limitation is that no other sessions can be connected to the source database while it is being copied.
CREATE DATABASE will fail if any other connection exists when it starts; during the copy operation,
new connections to the source database are prevented.

Two useful flags exist in pg_database for each database: the columns datistemplate
and datallowconn. datistemplate can be set to indicate that a database is intended as a
template for CREATE DATABASE. If this flag is set, the database can be cloned by any user with
CREATEDB privileges; if it is not set, only superusers and the owner of the database can clone it. If
datallowconn is false, then no new connections to that database will be allowed (but existing
sessions are not terminated simply by setting the flag false). The template0 database is normally
marked datallowconn = false to prevent its modification. Both template0 and template1
should always be marked with datistemplate = true.

Note: template1 and template0 do not have any special status beyond the fact that the name
template1 is the default source database name for CREATE DATABASE. For example, one could
drop template1 and recreate it from template0 without any ill effects. This course of action might
be advisable if one has carelessly added a bunch of junk in template1. (To delete template1, it
must have pg_database.datistemplate = false.)

The postgres database is also created when a database cluster is initialized. This database
is meant as a default database for users and applications to connect to. It is simply a copy of
template1 and can be dropped and recreated if necessary.

## 22.4. Database Configuration

Recall from Chapter 19 that the PostgreSQL server provides a large number of run-time configuration
variables. You can set database-specific default values for many of these settings.

For example, if for some reason you want to disable the GEQO optimizer for a given database, you’d
ordinarily have to either disable it for all databases or make sure that every connecting client is careful
to issue SET geqo TO off. To make this setting the default within a particular database, you can
execute the command:
ALTER DATABASE mydb SET geqo TO off;

This will save the setting (but not set it immediately). In subsequent connections to this database it will appear as though SET geqo TO off; had been executed just before the session started. Note that users can still alter this setting during their sessions; it will only be the default. To undo any such setting, use ALTER DATABASE dbname RESET varname.

22.5. Destroying a Database

Databases are destroyed with the command DROP DATABASE:

DROP DATABASE name;

Only the owner of the database, or a superuser, can drop a database. Dropping a database removes all objects that were contained within the database. The destruction of a database cannot be undone.

You cannot execute the DROP DATABASE command while connected to the victim database. You can, however, be connected to any other database, including the template1 database. template1 would be the only option for dropping the last user database of a given cluster.

For convenience, there is also a shell program to drop databases, dropdb:

dropdb dbname

(Unlike createdb, it is not the default action to drop the database with the current user name.)

22.6. Tablespaces

Tablespaces in PostgreSQL allow database administrators to define locations in the file system where the files representing database objects can be stored. Once created, a tablespace can be referred to by name when creating database objects.

By using tablespaces, an administrator can control the disk layout of a PostgreSQL installation. This is useful in at least two ways. First, if the partition or volume on which the cluster was initialized runs out of space and cannot be extended, a tablespace can be created on a different partition and used until the system can be reconfigured.

Second, tablespaces allow an administrator to use knowledge of the usage pattern of database objects to optimize performance. For example, an index which is very heavily used can be placed on a very fast, highly available disk, such as an expensive solid state device. At the same time a table storing archived data which is rarely used or not performance critical could be stored on a less expensive, slower disk system.

Warning

Even though located outside the main PostgreSQL data directory, tablespaces are an integral part of the database cluster and cannot be treated as an autonomous collection of data files. They are dependent on metadata contained in the main data directory, and therefore cannot be attached to a different database cluster or backed up individually. Similarly, if you lose a tablespace (file deletion, disk failure, etc), the database cluster might become unreadable or unable to start. Placing a tablespace on a temporary file system like a RAM disk risks the reliability of the entire cluster.
To define a tablespace, use the CREATE TABLESPACE command, for example::

```
CREATE TABLESPACE fastspace LOCATION '/ssd1/postgresql/data';
```

The location must be an existing, empty directory that is owned by the PostgreSQL operating system user. All objects subsequently created within the tablespace will be stored in files underneath this directory. The location must not be on removable or transient storage, as the cluster might fail to function if the tablespace is missing or lost.

**Note:** There is usually not much point in making more than one tablespace per logical file system, since you cannot control the location of individual files within a logical file system. However, PostgreSQL does not enforce any such limitation, and indeed it is not directly aware of the file system boundaries on your system. It just stores files in the directories you tell it to use.

Creation of the tablespace itself must be done as a database superuser, but after that you can allow ordinary database users to use it. To do that, grant them the `CREATE` privilege on it.

Tables, indexes, and entire databases can be assigned to particular tablespaces. To do so, a user with the `CREATE` privilege on a given tablespace must pass the tablespace name as a parameter to the relevant command. For example, the following creates a table in the tablespace `space1`:

```
CREATE TABLE foo(i int) TABLESPACE space1;
```

Alternatively, use the `default_tablespace` parameter:

```
SET default_tablespace = space1;
CREATE TABLE foo(i int);
```

When `default_tablespace` is set to anything but an empty string, it supplies an implicit `TABLESPACE` clause for `CREATE TABLE` and `CREATE INDEX` commands that do not have an explicit one.

There is also a `temp_tablespaces` parameter, which determines the placement of temporary tables and indexes, as well as temporary files that are used for purposes such as sorting large data sets. This can be a list of tablespace names, rather than only one, so that the load associated with temporary objects can be spread over multiple tablespaces. A random member of the list is picked each time a temporary object is to be created.

The tablespace associated with a database is used to store the system catalogs of that database. Furthermore, it is the default tablespace used for tables, indexes, and temporary files created within the database, if no `TABLESPACE` clause is given and no other selection is specified by `default_tablespace` or `temp_tablespaces` (as appropriate). If a database is created without specifying a tablespace for it, it uses the same tablespace as the template database it is copied from.

Two tablespaces are automatically created when the database cluster is initialized. The `pg_global` tablespace is used for shared system catalogs. The `pg_default` tablespace is the default tablespace of the `template1` and `template0` databases (and, therefore, will be the default tablespace for other databases as well, unless overridden by a `TABLESPACE` clause in `CREATE DATABASE`).

Once created, a tablescape can be used from any database, provided the requesting user has sufficient privilege. This means that a tablespace cannot be dropped until all objects in all databases using the tablespace have been removed.

To remove an empty tablespace, use the DROP TABLESPACE command.
To determine the set of existing tablespaces, examine the `pg_tablespace` system catalog, for example:

```
SELECT spcname FROM pg_tablespace;
```

The `psql` program’s `\db` meta-command is also useful for listing the existing tablespaces. PostgreSQL makes use of symbolic links to simplify the implementation of tablespaces. This means that tablespaces can be used only on systems that support symbolic links.

The directory `$PGDATA/pg_tblspc` contains symbolic links that point to each of the non-built-in tablespaces defined in the cluster. Although not recommended, it is possible to adjust the tablespace layout by hand by redefining these links. Under no circumstances perform this operation while the server is running. Note that in PostgreSQL 9.1 and earlier you will also need to update the `pg_tablespace` catalog with the new locations. (If you do not, `pg_dump` will continue to output the old tablespace locations.)
Chapter 23. Localization

This chapter describes the available localization features from the point of view of the administrator. PostgreSQL supports two localization facilities:

- Using the locale features of the operating system to provide locale-specific collation order, number formatting, translated messages, and other aspects. This is covered in Section 23.1 and Section 23.2.
- Providing a number of different character sets to support storing text in all kinds of languages, and providing character set translation between client and server. This is covered in Section 23.3.

23.1. Locale Support

Locale support refers to an application respecting cultural preferences regarding alphabets, sorting, number formatting, etc. PostgreSQL uses the standard ISO C and POSIX locale facilities provided by the server operating system. For additional information refer to the documentation of your system.

23.1.1. Overview

Locale support is automatically initialized when a database cluster is created using `initdb`. `initdb` will initialize the database cluster with the locale setting of its execution environment by default, so if your system is already set to use the locale that you want in your database cluster then there is nothing else you need to do. If you want to use a different locale (or you are not sure which locale your system is set to), you can instruct `initdb` exactly which locale to use by specifying the `--locale` option. For example:

```
initdb --locale=sv_SE
```

This example for Unix systems sets the locale to Swedish (sv) as spoken in Sweden (SE). Other possibilities might include `en_US` (U.S. English) and `fr_CA` (French Canadian). If more than one character set can be used for a locale then the specifications can take the form `language_territory.codeset`. For example, `fr_BE.UTF-8` represents the French language (fr) as spoken in Belgium (BE), with a UTF-8 character set encoding.

What locales are available on your system under what names depends on what was provided by the operating system vendor and what was installed. On most Unix systems, the command `locale -a` will provide a list of available locales. Windows uses more verbose locale names, such as `German_Germany` or `Swedish_Sweden.1252`, but the principles are the same.

Occasionally it is useful to mix rules from several locales, e.g., use English collation rules but Spanish messages. To support that, a set of locale subcategories exist that control only certain aspects of the localization rules:

<table>
<thead>
<tr>
<th><strong>LC_COLLATE</strong></th>
<th>String sort order</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>LC_CTYPE</strong></td>
<td>Character classification (What is a letter? Its upper-case equivalent?)</td>
</tr>
</tbody>
</table>
Chapter 23. Localization

<table>
<thead>
<tr>
<th>Category</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>LC_MESSAGES</td>
<td>Language of messages</td>
</tr>
<tr>
<td>LC_MONETARY</td>
<td>Formatting of currency amounts</td>
</tr>
<tr>
<td>LC_NUMERIC</td>
<td>Formatting of numbers</td>
</tr>
<tr>
<td>LC_TIME</td>
<td>Formatting of dates and times</td>
</tr>
</tbody>
</table>

The category names translate into names of initdb options to override the locale choice for a specific category. For instance, to set the locale to French Canadian, but use U.S. rules for formatting currency, use `initdb --locale=fr_CA --lc-monetary=en_US`.

If you want the system to behave as if it had no locale support, use the special locale name `C`, or equivalently `POSIX`.

Some locale categories must have their values fixed when the database is created. You can use different settings for different databases, but once a database is created, you cannot change them for that database anymore. `LC_COLLATE` and `LC_CTYPE` are these categories. They affect the sort order of indexes, so they must be kept fixed, or indexes on text columns would become corrupt. (But you can alleviate this restriction using collations, as discussed in Section 23.2.) The default values for these categories are determined when initdb is run, and those values are used when new databases are created, unless specified otherwise in the CREATE DATABASE command.

The other locale categories can be changed whenever desired by setting the server configuration parameters that have the same name as the locale categories (see Section 19.11.2 for details). The values that are chosen by initdb are actually only written into the configuration file `postgresql.conf` to serve as defaults when the server is started. If you remove these assignments from `postgresql.conf` then the server will inherit the settings from its execution environment.

Note that the locale behavior of the server is determined by the environment variables seen by the server, not by the environment of any client. Therefore, be careful to configure the correct locale settings before starting the server. A consequence of this is that if client and server are set up in different locales, messages might appear in different languages depending on where they originated.

**Note:** When we speak of inheriting the locale from the execution environment, this means the following on most operating systems: For a given locale category, say the collation, the following environment variables are consulted in this order until one is found to be set: `LC_ALL`, `LC_COLLATE` (or the variable corresponding to the respective category), `LANG`. If none of these environment variables are set then the locale defaults to `C`.

Some message localization libraries also look at the environment variable `LANGUAGE` which overrides all other locale settings for the purpose of setting the language of messages. If in doubt, please refer to the documentation of your operating system, in particular the documentation about gettext.

To enable messages to be translated to the user’s preferred language, NLS must have been selected at build time (`configure --enable-nls`). All other locale support is built in automatically.

### 23.1.2. Behavior

The locale settings influence the following SQL features:

- Sort order in queries using `ORDER BY` or the standard comparison operators on textual data
- The `upper`, `lower`, and `initcap` functions
Chapter 23. Localization

• Pattern matching operators (LIKE, SIMILAR TO, and POSIX-style regular expressions); locales affect both case insensitive matching and the classification of characters by character-class regular expressions
• The to_char family of functions
• The ability to use indexes with LIKE clauses

The drawback of using locales other than C or POSIX in PostgreSQL is its performance impact. It slows character handling and prevents ordinary indexes from being used by LIKE. For this reason use locales only if you actually need them.

As a workaround to allow PostgreSQL to use indexes with LIKE clauses under a non-C locale, several custom operator classes exist. These allow the creation of an index that performs a strict character-by-character comparison, ignoring locale comparison rules. Refer to Section 11.9 for more information. Another approach is to create indexes using the C collation, as discussed in Section 23.2.

23.1.3. Problems

If locale support doesn’t work according to the explanation above, check that the locale support in your operating system is correctly configured. To check what locales are installed on your system, you can use the command `locale -a` if your operating system provides it.

Check that PostgreSQL is actually using the locale that you think it is. The LC_COLLATE and LC_CTYPE settings are determined when a database is created, and cannot be changed except by creating a new database. Other locale settings including LC_MESSAGES and LC_MONETARY are initially determined by the environment the server is started in, but can be changed on-the-fly. You can check the active locale settings using the SHOW command.

The directory src/test/locale in the source distribution contains a test suite for PostgreSQL’s locale support.

Client applications that handle server-side errors by parsing the text of the error message will obviously have problems when the server’s messages are in a different language. Authors of such applications are advised to make use of the error code scheme instead.

Maintaining catalogs of message translations requires the on-going efforts of many volunteers that want to see PostgreSQL speak their preferred language well. If messages in your language are currently not available or not fully translated, your assistance would be appreciated. If you want to help, refer to Chapter 53 or write to the developers’ mailing list.

23.2. Collation Support

The collation feature allows specifying the sort order and character classification behavior of data per-column, or even per-operation. This alleviates the restriction that the LC_COLLATE and LC_CTYPE settings of a database cannot be changed after its creation.

23.2.1. Concepts

Conceptually, every expression of a collatable data type has a collation. (The built-in collatable data types are text, varchar, and char. User-defined base types can also be marked collatable, and of
course a domain over a collatable data type is collatable.) If the expression is a column reference, the
collation of the expression is the defined collation of the column. If the expression is a constant, the
collation is the default collation of the data type of the constant. The collation of a more complex
expression is derived from the collations of its inputs, as described below.

The collation of an expression can be the “default” collation, which means the locale settings defined
for the database. It is also possible for an expression’s collation to be indeterminate. In such cases,
ordering operations and other operations that need to know the collation will fail.

When the database system has to perform an ordering or a character classification, it uses the collation
of the input expression. This happens, for example, with ORDER BY clauses and function or operator
calls such as <. The collation to apply for an ORDER BY clause is simply the collation of the sort
key. The collation to apply for a function or operator call is derived from the arguments, as described
below. In addition to comparison operators, collations are taken into account by functions that convert
between lower and upper case letters, such as lower, upper, and initcap; by pattern matching
operators; and by to_char and related functions.

For a function or operator call, the collation that is derived by examining the argument collations is
used at run time for performing the specified operation. If the result of the function or operator call
is of a collatable data type, the collation is also used at parse time as the defined collation of the
function or operator expression, in case there is a surrounding expression that requires knowledge of
its collation.

The collation derivation of an expression can be implicit or explicit. This distinction affects how col-
lations are combined when multiple different collations appear in an expression. An explicit collation
derivation occurs when a COLLATE clause is used; all other collation derivations are implicit. When
multiple collations need to be combined, for example in a function call, the following rules are used:

1. If any input expression has an explicit collation derivation, then all explicitly derived collations
among the input expressions must be the same, otherwise an error is raised. If any explicitly
derived collation is present, that is the result of the collation combination.

2. Otherwise, all input expressions must have the same implicit collation derivation or the default
collation. If any non-default collation is present, that is the result of the collation combination.
   Otherwise, the result is the default collation.

3. If there are conflicting non-default implicit collations among the input expressions, then the com-
bination is deemed to have indeterminate collation. This is not an error condition unless the
particular function being invoked requires knowledge of the collation it should apply. If it does,
an error will be raised at run-time.

For example, consider this table definition:

```sql
CREATE TABLE test1 (  
  a text COLLATE "de_DE",  
  b text COLLATE "es_ES",  
  ...  
);
```

Then in

```sql
SELECT a < 'foo' FROM test1;
```

the < comparison is performed according to de_DE rules, because the expression combines an im-
plicitly derived collation with the default collation. But in

```sql
SELECT a < ('foo' COLLATE "fr_FR") FROM test1;
```
the comparison is performed using fr_FR rules, because the explicit collation derivation overrides the implicit one. Furthermore, given

```sql
SELECT a < b FROM test1;
```

the parser cannot determine which collation to apply, since the a and b columns have conflicting implicit collations. Since the < operator does need to know which collation to use, this will result in an error. The error can be resolved by attaching an explicit collation specifier to either input expression, thus:

```sql
SELECT a < b COLLATE "de_DE" FROM test1;
```

or equivalently

```sql
SELECT a COLLATE "de_DE" < b FROM test1;
```

On the other hand, the structurally similar case

```sql
SELECT a || b FROM test1;
```

does not result in an error, because the || operator does not care about collations: its result is the same regardless of the collation.

The collation assigned to a function or operator’s combined input expressions is also considered to apply to the function or operator’s result, if the function or operator delivers a result of a collatable data type. So, in

```sql
SELECT * FROM test1 ORDER BY a || 'foo';
```

the ordering will be done according to de_DE rules. But this query:

```sql
SELECT * FROM test1 ORDER BY a || b;
```

results in an error, because even though the || operator doesn’t need to know a collation, the ORDER BY clause does. As before, the conflict can be resolved with an explicit collation specifier:

```sql
SELECT * FROM test1 ORDER BY a || b COLLATE "fr_FR";
```

### 23.2.2. Managing Collations

A collation is an SQL schema object that maps an SQL name to operating system locales. In particular, it maps to a combination of LC_COLLATE and LC_CTYPE. (As the name would suggest, the main purpose of a collation is to set LC_COLLATE, which controls the sort order. But it is rarely necessary in practice to have an LC_CTYPE setting that is different from LC_COLLATE, so it is more convenient to collect these under one concept than to create another infrastructure for setting LC_CTYPE per expression.) Also, a collation is tied to a character set encoding (see Section 23.3). The same collation name may exist for different encodings.

On all platforms, the collations named default, C, and POSIX are available. Additional collations may be available depending on operating system support. The default collation selects the LC_COLLATE and LC_CTYPE values specified at database creation time. The C and POSIX collations both specify “traditional C” behavior, in which only the ASCII letters “A” through “Z” are treated as letters, and sorting is done strictly by character code byte values.
If the operating system provides support for using multiple locales within a single program (newlocale and related functions), then when a database cluster is initialized, initdb populates the system catalog pg_collation with collations based on all the locales it finds on the operating system at the time. For example, the operating system might provide a locale named de_DE.utf8. initdb would then create a collation named de_DE.utf8 for encoding UTF8 that has both LC_COLLATE and LC_CTYPE set to de_DE.utf8. It will also create a collation with the .utf8 tag stripped off the name. So you could also use the collation under the name de_DE, which is less cumbersome to write and makes the name less encoding-dependent. Note that, nevertheless, the initial set of collation names is platform-dependent.

In case a collation is needed that has different values for LC_COLLATE and LC_CTYPE, a new collation may be created using the CREATE COLLATION command. That command can also be used to create a new collation from an existing collation, which can be useful to be able to use operating-system-independent collation names in applications.

Within any particular database, only collations that use that database’s encoding are of interest. Other entries in pg_collation are ignored. Thus, a stripped collation name such as de_DE can be considered unique within a given database even though it would not be unique globally. Use of the stripped collation names is recommended, since it will make one less thing you need to change if you decide to change to another database encoding. Note however that the default, C, and POSIX collations can be used regardless of the database encoding.

PostgreSQL considers distinct collation objects to be incompatible even when they have identical properties. Thus for example,

```
SELECT a COLLATE "C" < b COLLATE "POSIX" FROM test1;
```

will draw an error even though the C and POSIX collations have identical behaviors. Mixing stripped and non-stripped collation names is therefore not recommended.

### 23.3. Character Set Support

The character set support in PostgreSQL allows you to store text in a variety of character sets (also called encodings), including single-byte character sets such as the ISO 8859 series and multiple-byte character sets such as EUC (Extended Unix Code), UTF-8, and Mule internal code. All supported character sets can be used transparently by clients, but a few are not supported for use within the server (that is, as a server-side encoding). The default character set is selected while initializing your PostgreSQL database cluster using initdb. It can be overridden when you create a database, so you can have multiple databases each with a different character set.

An important restriction, however, is that each database’s character set must be compatible with the database’s LC_CTYPE (character classification) and LC_COLLATE (string sort order) locale settings. For C or POSIX locale, any character set is allowed, but for other locales there is only one character set that will work correctly. (On Windows, however, UTF-8 encoding can be used with any locale.)

#### 23.3.1. Supported Character Sets

Table 23-1 shows the character sets available for use in PostgreSQL.

<table>
<thead>
<tr>
<th>Table 23-1. PostgreSQL Character Sets</th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>Name</td>
<td>Description</td>
<td>Language</td>
</tr>
<tr>
<td>--------------------</td>
<td>--------------------------------------------</td>
<td>------------</td>
</tr>
<tr>
<td>BIG5</td>
<td>Big Five</td>
<td>Traditional Chinese</td>
</tr>
<tr>
<td>EUC_CN</td>
<td>Extended UNIX Code-CN</td>
<td>Simplified Chinese</td>
</tr>
<tr>
<td>EUC_TW</td>
<td>Extended UNIX Code-TW</td>
<td>Traditional Chinese, Taiwanese</td>
</tr>
<tr>
<td>GB18030</td>
<td>National Standard</td>
<td>Chinese</td>
</tr>
<tr>
<td>GBK</td>
<td>Extended National Standard</td>
<td>Simplified Chinese</td>
</tr>
<tr>
<td>ISO_8859_5</td>
<td>ISO 8859-5, ECMA 113</td>
<td>Latin/Cyrillic</td>
</tr>
<tr>
<td>ISO_8859_6</td>
<td>ISO 8859-6, ECMA 114</td>
<td>Latin/Arabic</td>
</tr>
<tr>
<td>ISO_8859_7</td>
<td>ISO 8859-7, ECMA 118</td>
<td>Latin/Greek</td>
</tr>
<tr>
<td>ISO_8859_8</td>
<td>ISO 8859-8, ECMA 121</td>
<td>Latin/Hebrew</td>
</tr>
<tr>
<td>JOHAB</td>
<td>JOHAB</td>
<td>Korean (Hangul)</td>
</tr>
<tr>
<td>KOI8R</td>
<td>KOI8-R</td>
<td>Cyrillic (Russian)</td>
</tr>
<tr>
<td>KOI8U</td>
<td>KOI8-U</td>
<td>Cyrillic (Ukrainian)</td>
</tr>
<tr>
<td>LATIN1</td>
<td>ISO 8859-1, ECMA 94</td>
<td>Western European</td>
</tr>
<tr>
<td>LATIN2</td>
<td>ISO 8859-2, ECMA 94</td>
<td>Central European</td>
</tr>
<tr>
<td>LATIN3</td>
<td>ISO 8859-3, ECMA 94</td>
<td>South European</td>
</tr>
<tr>
<td>LATIN4</td>
<td>ISO 8859-4, ECMA 94</td>
<td>North European</td>
</tr>
<tr>
<td>Name</td>
<td>Description</td>
<td>Language</td>
</tr>
<tr>
<td>------------</td>
<td>------------------------------</td>
<td>------------</td>
</tr>
<tr>
<td>LATIN5</td>
<td>ISO 8859-9, ECMA 128</td>
<td>Turkish</td>
</tr>
<tr>
<td>LATIN6</td>
<td>ISO 8859-10, ECMA 144</td>
<td>Nordic</td>
</tr>
<tr>
<td>LATIN7</td>
<td>ISO 8859-13</td>
<td>Baltic</td>
</tr>
<tr>
<td>LATIN8</td>
<td>ISO 8859-14</td>
<td>Celtic</td>
</tr>
<tr>
<td>LATIN9</td>
<td>ISO 8859-15</td>
<td>LATIN1 with Euro and accents</td>
</tr>
<tr>
<td>LATIN10</td>
<td>ISO 8859-16, ASRO SR 14111</td>
<td>Romanian</td>
</tr>
<tr>
<td>MULE_INTERNAL</td>
<td>Mule internal code</td>
<td>Multilingual Emacs</td>
</tr>
<tr>
<td>SJIS</td>
<td>Shift JIS</td>
<td>Japanese</td>
</tr>
<tr>
<td>SHIFT_JIS_2004</td>
<td>Shift JIS, JIS X 0213</td>
<td>Japanese</td>
</tr>
<tr>
<td>SQL_ASCII</td>
<td>unspecified (see text)</td>
<td>any</td>
</tr>
<tr>
<td>UHC</td>
<td>Unified Hangul Code</td>
<td>Korean</td>
</tr>
<tr>
<td>UTF8</td>
<td>Unicode, 8-bit all</td>
<td>all</td>
</tr>
<tr>
<td>WIN866</td>
<td>Windows CP866</td>
<td>Cyrillic</td>
</tr>
<tr>
<td>WIN874</td>
<td>Windows CP874</td>
<td>Thai</td>
</tr>
<tr>
<td>WIN1250</td>
<td>Windows CP1250</td>
<td>Central European</td>
</tr>
<tr>
<td>WIN1251</td>
<td>Windows CP1251</td>
<td>Cyrillic</td>
</tr>
<tr>
<td>WIN1252</td>
<td>Windows CP1252</td>
<td>Western European</td>
</tr>
<tr>
<td>WIN1253</td>
<td>Windows CP1253</td>
<td>Greek</td>
</tr>
<tr>
<td>WIN1254</td>
<td>Windows CP1254</td>
<td>Turkish</td>
</tr>
<tr>
<td>WIN1255</td>
<td>Windows CP1255</td>
<td>Hebrew</td>
</tr>
<tr>
<td>WIN1256</td>
<td>Windows CP1256</td>
<td>Arabic</td>
</tr>
<tr>
<td>WIN1257</td>
<td>Windows CP1257</td>
<td>Baltic</td>
</tr>
</tbody>
</table>
Chapter 23. Localization

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
<th>Language</th>
<th>Server?</th>
<th>Bytes/Char</th>
<th>Aliases</th>
</tr>
</thead>
<tbody>
<tr>
<td>WIN1258</td>
<td>Windows CP1258</td>
<td>Vietnamese</td>
<td>Yes</td>
<td>1</td>
<td>ABC, TCVN, TCVN5712, VSCII</td>
</tr>
</tbody>
</table>

Not all client APIs support all the listed character sets. For example, the PostgreSQL JDBC driver does not support `MULE_INTERNAL`, `LATIN6`, `LATIN8`, and `LATIN10`.

The `SQL_ASCII` setting behaves considerably differently from the other settings. When the server character set is `SQL_ASCII`, the server interprets byte values 0-127 according to the ASCII standard, while byte values 128-255 are taken as uninterpreted characters. No encoding conversion will be done when the setting is `SQL_ASCII`. Thus, this setting is not so much a declaration that a specific encoding is in use, as a declaration of ignorance about the encoding. In most cases, if you are working with any non-ASCII data, it is unwise to use the `SQL_ASCII` setting because PostgreSQL will be unable to help you by converting or validating non-ASCII characters.

### 23.3.2. Setting the Character Set

`initdb` defines the default character set (encoding) for a PostgreSQL cluster. For example,

```
initdb -E EUC_JP
```

sets the default character set to `EUC_JP` (Extended Unix Code for Japanese). You can use `--encoding` instead of `-E` if you prefer longer option strings. If no `-E` or `--encoding` option is given, `initdb` attempts to determine the appropriate encoding to use based on the specified or default locale.

You can specify a non-default encoding at database creation time, provided that the encoding is compatible with the selected locale:

```
createdb -E EUC_KR -T template0 --lc-collate=ko_KR.euckr --lc-ctype=ko_KR.euckr korean
```

This will create a database named `korean` that uses the character set `EUC_KR`, and locale `ko_KR`. Another way to accomplish this is to use this SQL command:

```
CREATE DATABASE korean WITH ENCODING 'EUC_KR' LC_COLLATE='ko_KR.euckr' LC_CTYPE='ko_KR.euckr' TEMPLATE=template0;
```

Notice that the above commands specify copying the `template0` database. When copying any other database, the encoding and locale settings cannot be changed from those of the source database, because that might result in corrupt data. For more information see Section 22.3.

The encoding for a database is stored in the system catalog `pg_database`. You can see it by using the `psql -l` option or the `\l` command.

```
$ psql -l
```

```
List of databases
Name           | Owner | Encoding | Collation | Ctype | Access Privileges
--------------|-------|----------|-----------|-------|------------------
clocaledb      | hlinnaka | SQL_ASCII | C | C |
englishdb      | hlinnaka | UTF8     | en_GB.UTF8 | en_GB.UTF8 |
japanese       | hlinnaka | UTF8     | ja_JP.UTF8 | ja_JP.UTF8 |
korean         | hlinnaka | EUC_KR   | ko_KR.euckr | ko_KR.euckr |
postgres       | hlinnaka | UTF8     | fi_FI.UTF8 | fi_FI.UTF8 |
template0      | hlinnaka | UTF8     | fi_FI.UTF8 | fi_FI.UTF8 | {=c/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnaka,hlinnaka=CTc/hlinnakan |
```
Important: On most modern operating systems, PostgreSQL can determine which character set is implied by the \texttt{LC\_CTYPE} setting, and it will enforce that only the matching database encoding is used. On older systems it is your responsibility to ensure that you use the encoding expected by the locale you have selected. A mistake in this area is likely to lead to strange behavior of locale-dependent operations such as sorting.

PostgreSQL will allow superusers to create databases with \texttt{SQL\_ASCII} encoding even when \texttt{LC\_CTYPE} is not \texttt{C} or \texttt{POSIX}. As noted above, \texttt{SQL\_ASCII} does not enforce that the data stored in the database has any particular encoding, and so this choice poses risks of locale-dependent misbehavior. Using this combination of settings is deprecated and may someday be forbidden altogether.

### 23.3.3. Automatic Character Set Conversion Between Server and Client

PostgreSQL supports automatic character set conversion between server and client for certain character set combinations. The conversion information is stored in the \texttt{pg\_conversion} system catalog. PostgreSQL comes with some predefined conversions, as shown in Table 23-2. You can create a new conversion using the SQL command \texttt{CREATE CONVERSION}.

**Table 23-2. Client/Server Character Set Conversions**

<table>
<thead>
<tr>
<th>Server Character Set</th>
<th>Available Client Character Sets</th>
</tr>
</thead>
<tbody>
<tr>
<td>BIG5</td>
<td>\textit{not supported as a server encoding}</td>
</tr>
<tr>
<td>EUC_CN</td>
<td>EUC_CN, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>EUC_JP</td>
<td>EUC_JP, MULE_INTERNAL, SJIS, UTF8</td>
</tr>
<tr>
<td>EUC_JIS_2004</td>
<td>EUC_JIS_2004, SHIFT_JIS_2004, UTF8</td>
</tr>
<tr>
<td>EUC_KR</td>
<td>EUC_KR, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>EUC_TW</td>
<td>EUC_TW, BIG5, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>GB18030</td>
<td>\textit{not supported as a server encoding}</td>
</tr>
<tr>
<td>GBK</td>
<td>\textit{not supported as a server encoding}</td>
</tr>
<tr>
<td>ISO_8859_5</td>
<td>ISO_8859_5, KOI8R, MULE_INTERNAL, UTF8, WIN866, WIN1251</td>
</tr>
<tr>
<td>ISO_8859_6</td>
<td>ISO_8859_6, UTF8</td>
</tr>
<tr>
<td>ISO_8859_7</td>
<td>ISO_8859_7, UTF8</td>
</tr>
<tr>
<td>ISO_8859_8</td>
<td>ISO_8859_8, UTF8</td>
</tr>
<tr>
<td>JOHAB</td>
<td>\textit{not supported as a server encoding}</td>
</tr>
<tr>
<td>KOI8R</td>
<td>KOI8R, ISO_8859_5, MULE_INTERNAL, UTF8, WIN866, WIN1251</td>
</tr>
<tr>
<td>KOI8U</td>
<td>KOI8U, UTF8</td>
</tr>
<tr>
<td>LATIN1</td>
<td>LATIN1, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>LATIN2</td>
<td>LATIN2, MULE_INTERNAL, UTF8, WIN1250</td>
</tr>
</tbody>
</table>
Chapter 23. Localization

### Server Character Set

<table>
<thead>
<tr>
<th>Server Character Set</th>
<th>Available Client Character Sets</th>
</tr>
</thead>
<tbody>
<tr>
<td>LATIN3</td>
<td>LATIN3, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>LATIN4</td>
<td>LATIN4, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>LATIN5</td>
<td>LATIN5, UTF8</td>
</tr>
<tr>
<td>LATIN6</td>
<td>LATIN6, UTF8</td>
</tr>
<tr>
<td>LATIN7</td>
<td>LATIN7, UTF8</td>
</tr>
<tr>
<td>LATIN8</td>
<td>LATIN8, UTF8</td>
</tr>
<tr>
<td>LATIN9</td>
<td>LATIN9, UTF8</td>
</tr>
<tr>
<td>LATIN10</td>
<td>LATIN10, UTF8</td>
</tr>
<tr>
<td>MULE_INTERNAL</td>
<td>MULE_INTERNAL, BIG5, EUC_CN, EUC_JP, EUC_TW, ISO_8859_5, KOI8R, LATIN1 to LATIN4, SJIS, WIN866, WIN1250, WIN1251</td>
</tr>
<tr>
<td>SJIS</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>SHIFT_JIS_2004</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>SQL_ASCII</td>
<td>any (no conversion will be performed)</td>
</tr>
<tr>
<td>UHC</td>
<td>not supported as a server encoding</td>
</tr>
<tr>
<td>UTF8</td>
<td>all supported encodings</td>
</tr>
<tr>
<td>WIN866</td>
<td>WIN866, ISO_8859_5, KOI8R, MULE_INTERNAL, UTF8, WIN1251</td>
</tr>
<tr>
<td>WIN874</td>
<td>WIN874, UTF8</td>
</tr>
<tr>
<td>WIN1250</td>
<td>WIN1250, LATIN2, MULE_INTERNAL, UTF8</td>
</tr>
<tr>
<td>WIN1251</td>
<td>WIN1251, ISO_8859_5, KOI8R, MULE_INTERNAL, UTF8, WIN866</td>
</tr>
<tr>
<td>WIN1252</td>
<td>WIN1252, UTF8</td>
</tr>
<tr>
<td>WIN1253</td>
<td>WIN1253, UTF8</td>
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<tr>
<td>WIN1254</td>
<td>WIN1254, UTF8</td>
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<td>WIN1255</td>
<td>WIN1255, UTF8</td>
</tr>
<tr>
<td>WIN1256</td>
<td>WIN1256, UTF8</td>
</tr>
<tr>
<td>WIN1257</td>
<td>WIN1257, UTF8</td>
</tr>
<tr>
<td>WIN1258</td>
<td>WIN1258, UTF8</td>
</tr>
</tbody>
</table>

To enable automatic character set conversion, you have to tell PostgreSQL the character set (encoding) you would like to use in the client. There are several ways to accomplish this:

- **Using the \encoding command in psql.** \encoding allows you to change client encoding on the fly. For example, to change the encoding to SJIS, type:

  \encoding SJIS

- **libpq (Section 32.10)** has functions to control the client encoding.

- **Using SET client_encoding TO.** Setting the client encoding can be done with this SQL command:

  ```sql
  SET CLIENT_ENCODING TO 'value';
  ```

  Also you can use the standard SQL syntax **SET NAMES** for this purpose:
SET NAMES 'value';
To query the current client encoding:
SHOW client_encoding;
To return to the default encoding:
RESET client_encoding;

- Using PGCLIENTENCODING. If the environment variable PGCLIENTENCODING is defined in the client’s environment, that client encoding is automatically selected when a connection to the server is made. (This can subsequently be overridden using any of the other methods mentioned above.)

- Using the configuration variable client_encoding. If the client_encoding variable is set, that client encoding is automatically selected when a connection to the server is made. (This can subsequently be overridden using any of the other methods mentioned above.)

If the conversion of a particular character is not possible — suppose you chose EUC_JP for the server and LATIN1 for the client, and some Japanese characters are returned that do not have a representation in LATIN1 — an error is reported.

If the client character set is defined as SQL_ASCII, encoding conversion is disabled, regardless of the server’s character set. Just as for the server, use of SQL_ASCII is unwise unless you are working with all-ASCII data.

23.3.4. Further Reading

These are good sources to start learning about various kinds of encoding systems.

CJKV Information Processing: Chinese, Japanese, Korean & Vietnamese Computing

Contains detailed explanations of EUC_JP, EUC_CN, EUC_KR, EUC_TW.

http://www.unicode.org/

The web site of the Unicode Consortium.

RFC 3629

UTF-8 (8-bit UCS/Unicode Transformation Format) is defined here.
Chapter 24. Routine Database Maintenance Tasks

PostgreSQL, like any database software, requires that certain tasks be performed regularly to achieve optimum performance. The tasks discussed here are required, but they are repetitive in nature and can easily be automated using standard tools such as cron scripts or Windows’ Task Scheduler. It is the database administrator’s responsibility to set up appropriate scripts, and to check that they execute successfully.

One obvious maintenance task is the creation of backup copies of the data on a regular schedule. Without a recent backup, you have no chance of recovery after a catastrophe (disk failure, fire, mistakenly dropping a critical table, etc.). The backup and recovery mechanisms available in PostgreSQL are discussed at length in Chapter 25.

The other main category of maintenance task is periodic “vacuuming” of the database. This activity is discussed in Section 24.1. Closely related to this is updating the statistics that will be used by the query planner, as discussed in Section 24.1.3.

Another task that might need periodic attention is log file management. This is discussed in Section 24.3.

check_postgres1 is available for monitoring database health and reporting unusual conditions. check_postgres integrates with Nagios and MRTG, but can be run standalone too.

PostgreSQL is low-maintenance compared to some other database management systems. Nonetheless, appropriate attention to these tasks will go far towards ensuring a pleasant and productive experience with the system.

24.1. Routine Vacuuming

PostgreSQL databases require periodic maintenance known as vacuuming. For many installations, it is sufficient to let vacuuming be performed by the autovacuum daemon, which is described in Section 24.1.6. You might need to adjust the autovacuuming parameters described there to obtain best results for your situation. Some database administrators will want to supplement or replace the daemon’s activities with manually-managed VACUUM commands, which typically are executed according to a schedule by cron or Task Scheduler scripts. To set up manually-managed vacuuming properly, it is essential to understand the issues discussed in the next few subsections. Administrators who rely on autovacuuming may still wish to skim this material to help them understand and adjust autovacuuming.

24.1.1. Vacuuming Basics

PostgreSQL’s VACUUM command has to process each table on a regular basis for several reasons:

1. To recover or reuse disk space occupied by updated or deleted rows.
2. To update data statistics used by the PostgreSQL query planner.
3. To update the visibility map, which speeds up index-only scans.

1. https://bucardo.org/check_postgres/
4. To protect against loss of very old data due to transaction ID wraparound or multixact ID wraparound.

Each of these reasons dictates performing VACUUM operations of varying frequency and scope, as explained in the following subsections.

There are two variants of VACUUM: standard VACUUM and VACUUM FULL. VACUUM FULL can reclaim more disk space but runs much more slowly. Also, the standard form of VACUUM can run in parallel with production database operations. (Commands such as SELECT, INSERT, UPDATE, and DELETE will continue to function normally, though you will not be able to modify the definition of a table with commands such as ALTER TABLE while it is being vacuumed.) VACUUM FULL requires exclusive lock on the table it is working on, and therefore cannot be done in parallel with other use of the table. Generally, therefore, administrators should strive to use standard VACUUM and avoid VACUUM FULL.

VACUUM creates a substantial amount of I/O traffic, which can cause poor performance for other active sessions. There are configuration parameters that can be adjusted to reduce the performance impact of background vacuuming — see Section 19.4.4.

### 24.1.2. Recovering Disk Space

In PostgreSQL, an UPDATE or DELETE of a row does not immediately remove the old version of the row. This approach is necessary to gain the benefits of multiversion concurrency control (MVCC, see Chapter 13): the row version must not be deleted while it is still potentially visible to other transactions. But eventually, an outdated or deleted row version is no longer of interest to any transaction. The space it occupies must then be reclaimed for reuse by new rows, to avoid unbounded growth of disk space requirements. This is done by running VACUUM.

The standard form of VACUUM removes dead row versions in tables and indexes and marks the space available for future reuse. However, it will not return the space to the operating system, except in the special case where one or more pages at the end of a table become entirely free and an exclusive table lock can be easily obtained. In contrast, VACUUM FULL actively compacts tables by writing a complete new version of the table file with no dead space. This minimizes the size of the table, but can take a long time. It also requires extra disk space for the new copy of the table, until the operation completes.

The usual goal of routine vacuuming is to do standard VACUUMs often enough to avoid needing VACUUM FULL. The autovacuum daemon attempts to work this way, and in fact will never issue VACUUM FULL. In this approach, the idea is not to keep tables at their minimum size, but to maintain steady-state usage of disk space: each table occupies space equivalent to its minimum size plus however much space gets used up between vacuumings. Although VACUUM FULL can be used to shrink a table back to its minimum size and return the disk space to the operating system, there is not much point in this if the table will just grow again in the future. Thus, moderately-frequent standard VACUUM runs are a better approach than infrequent VACUUM FULL runs for maintaining heavily-updated tables.

Some administrators prefer to schedule vacuuming themselves, for example doing all the work at night when load is low. The difficulty with doing vacuuming according to a fixed schedule is that if a table has an unexpected spike in update activity, it may get bloated to the point that VACUUM FULL is really necessary to reclaim space. Using the autovacuum daemon alleviates this problem, since the daemon schedules vacuuming dynamically in response to update activity. It is unwise to disable the daemon completely unless you have an extremely predictable workload. One possible compromise is to set the daemon’s parameters so that it will only react to unusually heavy update activity, thus keeping things from getting out of hand, while scheduled VACUUMs are expected to do the bulk of the work when the load is typical.
Chapter 24. Routine Database Maintenance Tasks

For those not using autovacuum, a typical approach is to schedule a database-wide \texttt{VACUUM} once a day during a low-usage period, supplemented by more frequent vacuuming of heavily-updated tables as necessary. (Some installations with extremely high update rates vacuum their busiest tables as often as once every few minutes.) If you have multiple databases in a cluster, don’t forget to \texttt{VACUUM} each one; the program vacuumdb might be helpful.

**Tip:** Plain \texttt{VACUUM} may not be satisfactory when a table contains large numbers of dead row versions as a result of massive update or delete activity. If you have such a table and you need to reclaim the excess disk space it occupies, you will need to use \texttt{VACUUM FULL}, or alternatively \texttt{CLUSTER} or one of the table-rewriting variants of ALTER TABLE. These commands rewrite an entire new copy of the table and build new indexes for it. All these options require exclusive lock. Note that they also temporarily use extra disk space approximately equal to the size of the table, since the old copies of the table and indexes can’t be released until the new ones are complete.

**Tip:** If you have a table whose entire contents are deleted on a periodic basis, consider doing it with \texttt{TRUNCATE} rather than using \texttt{DELETE} followed by \texttt{VACUUM}. \texttt{TRUNCATE} removes the entire content of the table immediately, without requiring a subsequent \texttt{VACUUM} or \texttt{VACUUM FULL} to reclaim the now-unused disk space. The disadvantage is that strict MVCC semantics are violated.

### 24.1.3. Updating Planner Statistics

The PostgreSQL query planner relies on statistical information about the contents of tables in order to generate good plans for queries. These statistics are gathered by the \texttt{ANALYZE} command, which can be invoked by itself or as an optional step in \texttt{VACUUM}. It is important to have reasonably accurate statistics, otherwise poor choices of plans might degrade database performance.

The autovacuum daemon, if enabled, will automatically issue \texttt{ANALYZE} commands whenever the content of a table has changed sufficiently. However, administrators might prefer to rely on manually-scheduled \texttt{ANALYZE} operations, particularly if it is known that update activity on a table will not affect the statistics of “interesting” columns. The daemon schedules \texttt{ANALYZE} strictly as a function of the number of rows inserted or updated; it has no knowledge of whether that will lead to meaningful statistical changes.

As with vacuuming for space recovery, frequent updates of statistics are more useful for heavily-updated tables than for seldom-updated ones. But even for a heavily-updated table, there might be no need for statistics updates if the statistical distribution of the data is not changing much. A simple rule of thumb is to think about how much the minimum and maximum values of the columns in the table change. For example, a \texttt{timestamp} column that contains the time of row update will have a constantly-increasing maximum value as rows are added and updated; such a column will probably need more frequent statistics updates than, say, a column containing URLs for pages accessed on a website. The URL column might receive changes just as often, but the statistical distribution of its values probably changes relatively slowly.

It is possible to run \texttt{ANALYZE} on specific tables and even just specific columns of a table, so the flexibility exists to update some statistics more frequently than others if your application requires it. In practice, however, it is usually best to just analyze the entire database, because it is a fast operation. \texttt{ANALYZE} uses a statistically random sampling of the rows of a table rather than reading every single row.
Tip: Although per-column tweaking of `ANALYZE` frequency might not be very productive, you might find it worthwhile to do per-column adjustment of the level of detail of the statistics collected by `ANALYZE`. Columns that are heavily used in `WHERE` clauses and have highly irregular data distributions might require a finer-grain data histogram than other columns. See `ALTER TABLE SET STATISTICS`, or change the database-wide default using the `default_statistics_target` configuration parameter.

Also, by default there is limited information available about the selectivity of functions. However, if you create an expression index that uses a function call, useful statistics will be gathered about the function, which can greatly improve query plans that use the expression index.

Tip: The autovacuum daemon does not issue `ANALYZE` commands for foreign tables, since it has no means of determining how often that might be useful. If your queries require statistics on foreign tables for proper planning, it’s a good idea to run manually-managed `ANALYZE` commands on those tables on a suitable schedule.

### 24.1.4. Updating The Visibility Map

Vacuum maintains a visibility map for each table to keep track of which pages contain only tuples that are known to be visible to all active transactions (and all future transactions, until the page is again modified). This has two purposes. First, vacuum itself can skip such pages on the next run, since there is nothing to clean up.

Second, it allows PostgreSQL to answer some queries using only the index, without reference to the underlying table. Since PostgreSQL indexes don’t contain tuple visibility information, a normal index scan fetches the heap tuple for each matching index entry, to check whether it should be seen by the current transaction. An index-only scan, on the other hand, checks the visibility map first. If it’s known that all tuples on the page are visible, the heap fetch can be skipped. This is most useful on large data sets where the visibility map can prevent disk accesses. The visibility map is vastly smaller than the heap, so it can easily be cached even when the heap is very large.

### 24.1.5. Preventing Transaction ID Wraparound Failures

PostgreSQL’s MVCC transaction semantics depend on being able to compare transaction ID (XID) numbers: a row version with an insertion XID greater than the current transaction’s XID is “in the future” and should not be visible to the current transaction. But since transaction IDs have limited size (32 bits) a cluster that runs for a long time (more than 4 billion transactions) would suffer transaction ID wraparound: the XID counter wraps around to zero, and all of a sudden transactions that were in the past appear to be in the future — which means their output become invisible. In short, catastrophic data loss. (Actually the data is still there, but that’s cold comfort if you cannot get at it.) To avoid this, it is necessary to vacuum every table in every database at least once every two billion transactions.

The reason that periodic vacuuming solves the problem is that `VACUUM` will mark rows as frozen, indicating that they were inserted by a transaction that committed sufficiently far in the past that the effects of the inserting transaction are certain to be visible to all current and future transactions. Normal XIDs are compared using modulo-2\(^{32}\) arithmetic. This means that for every normal XID, there are two billion XIDs that are “older” and two billion that are “newer”; another way to say it is that the normal XID space is circular with no endpoint. Therefore, once a row version has been created with a particular normal XID, the row version will appear to be “in the past” for the next
two billion transactions, no matter which normal XID we are talking about. If the row version still exists after more than two billion transactions, it will suddenly appear to be in the future. To prevent this, PostgreSQL reserves a special XID, FrozenTransactionId, which does not follow the normal XID comparison rules and is always considered older than every normal XID. Frozen row versions are treated as if the inserting XID were FrozenTransactionId, so that they will appear to be “in the past” to all normal transactions regardless of wraparound issues, and so such row versions will be valid until deleted, no matter how long that is.

**Note:** In PostgreSQL versions before 9.4, freezing was implemented by actually replacing a row’s insertion XID with FrozenTransactionId, which was visible in the row’s xmin system column. Newer versions just set a flag bit, preserving the row’s original xmin for possible forensic use. However, rows with xmin equal to FrozenTransactionId (2) may still be found in databases pg_upgrade’d from pre-9.4 versions.

Also, system catalogs may contain rows with xmin equal to BootstrapTransactionId (1), indicating that they were inserted during the first phase of initdb. Like FrozenTransactionId, this special XID is treated as older than every normal XID.

vacuum_freeze_min_age controls how old an XID value has to be before rows bearing that XID will be frozen. Increasing this setting may avoid unnecessary work if the rows that would otherwise be frozen will soon be modified again, but decreasing this setting increases the number of transactions that can elapse before the table must be vacuumed again.

VACUUM uses the visibility map to determine which pages of a table must be scanned. Normally, it will skip pages that don’t have any dead row versions even if those pages might still have row versions with old XID values. Therefore, normal VACUUM won’t always freeze every old row version in the table. Periodically, VACUUM will perform an *aggressive vacuum*, skipping only those pages which contain neither dead rows nor any unfrozen XID or MXID values. vacuum_freeze_table_age controls when VACUUM does that: all-visible but not all-frozen pages are scanned if the number of transactions that have passed since the last such scan is greater than vacuum_freeze_table_age minus vacuum_freeze_min_age. Setting vacuum_freeze_table_age to 0 forces VACUUM to use this more aggressive strategy for all scans.

The maximum time that a table can go unvacuumed is two billion transactions minus the vacuum_freeze_min_age value at the time of the last aggressive vacuum. If it were to go unvacuumed for longer than that, data loss could result. To ensure that this does not happen, autovacuum is invoked on any table that might contain unfrozen rows with XIDs older than the age specified by the configuration parameter autovacuum_freeze_max_age. (This will happen even if autovacuum is disabled.)

This implies that if a table is not otherwise vacuumed, autovacuum will be invoked on it approximately once every autovacuum_freeze_max_age minus vacuum_freeze_min_age transactions. For tables that are regularly vacuumed for space reclamation purposes, this is of little importance. However, for static tables (including tables that receive inserts, but no updates or deletes), there is no need to vacuum for space reclamation, so it can be useful to try to maximize the interval between forced autovacuums on very large static tables. Obviously one can do this either by increasing autovacuum_freeze_max_age or decreasing vacuum_freeze_min_age.

The effective maximum for vacuum_freeze_table_age is 0.95 * autovacuum_freeze_max_age; a setting higher than that will be capped to the maximum. A value higher than autovacuum_freeze_max_age wouldn’t make sense because an anti-wraparound autovacuum would be triggered at that point anyway, and the 0.95 multiplier leaves some breathing room to run a manual VACUUM before that happens. As a rule of thumb, vacuum_freeze_table_age should be set to a value somewhat below
autovacuum_freeze_max_age, leaving enough gap so that a regularly scheduled VACUUM or an autovacuum triggered by normal delete and update activity is run in that window. Setting it too close could lead to anti-wraparound autovacuums, even though the table was recently vacuumed to reclaim space, whereas lower values lead to more frequent aggressive vacuuming.

The sole disadvantage of increasing autovacuum_freeze_max_age (and vacuum_freeze_table_age along with it) is that the pg_clog subdirectory of the database cluster will take more space, because it must store the commit status of all transactions back to the autovacuum_freeze_max_age horizon. The commit status uses two bits per transaction, so if autovacuum_freeze_max_age is set to its maximum allowed value of two billion, pg_clog can be expected to grow to about half a gigabyte. If this is trivial compared to your total database size, setting autovacuum_freeze_max_age to its maximum allowed value is recommended. Otherwise, set it depending on what you are willing to allow for pg_clog storage. (The default, 200 million transactions, translates to about 50MB of pg_clog storage.)

One disadvantage of decreasing vacuum_freeze_min_age is that it might cause VACUUM to do useless work: freezing a row version is a waste of time if the row is modified soon thereafter (causing it to acquire a new XID). So the setting should be large enough that rows are not frozen until they are unlikely to change any more.

To track the age of the oldest unfrozen XIDs in a database, VACUUM stores XID statistics in the system tables pg_class and pg_database. In particular, the relfrozenxid column of a table’s pg_class row contains the freeze cutoff XID that was used by the last aggressive VACUUM for that table. All rows inserted by transactions with XIDs older than this cutoff XID are guaranteed to have been frozen. Similarly, the datfrozenxid column of a database’s pg_database row is a lower bound on the unfrozen XIDs appearing in that database — it is just the minimum of the per-table relfrozenxid values within the database. A convenient way to examine this information is to execute queries such as:

```sql
SELECT c.oid::regclass as table_name,
       greatest(age(c.relfrozenxid),age(t.relfrozenxid)) as age
FROM pg_class c
LEFT JOIN pg_class t ON c.reltoastrelid = t.oid
WHERE c.relkind IN ('r', 'm');

SELECT datname, age(datfrozenxid) FROM pg_database;
```

The age column measures the number of transactions from the cutoff XID to the current transaction’s XID.

VACUUM normally only scans pages that have been modified since the last vacuum, but relfrozenxid can only be advanced when every page of the table that might contain unfrozen XIDs is scanned. This happens when relfrozenxid is more than vacuum_freeze_table_age transactions old, when VACUUM’s FREEZE option is used, or when all pages that are not already all-frozen happen to require vacuuming to remove dead row versions. When VACUUM scans every page in the table that is not already all-frozen, it should set age(relfrozenxid) to a value just a little more than the vacuum_freeze_min_age setting that was used (more by the number of transactions started since the VACUUM started). If no relfrozenxid-advancing VACUUM is issued on the table until autovacuum_freeze_max_age is reached, an autovacuum will soon be forced for the table.

If for some reason autovacuum fails to clear old XIDs from a table, the system will begin to emit warning messages like this when the database’s oldest XIDs reach ten million transactions from the wraparound point:

```
WARNING: database "mydb" must be vacuumed within 177009986 transactions
HINT: To avoid a database shutdown, execute a database-wide VACUUM in "mydb".
```
(A manual \texttt{VACUUM} should fix the problem, as suggested by the hint; but note that the \texttt{VACUUM} must be performed by a superuser, else it will fail to process system catalogs and thus not be able to advance the database's \texttt{datfrozenxid}.) If these warnings are ignored, the system will shut down and refuse to start any new transactions once there are fewer than 1 million transactions left until wraparound:

```
ERROR: database is not accepting commands to avoid wraparound data loss in database "mydb"
HINT: Stop the postmaster and vacuum that database in single-user mode.
```

The 1-million-transaction safety margin exists to let the administrator recover without data loss, by manually executing the required \texttt{VACUUM} commands. However, since the system will not execute commands once it has gone into the safety shutdown mode, the only way to do this is to stop the server and start the server in single-user mode to execute \texttt{VACUUM}. The shutdown mode is not enforced in single-user mode. See the postgres reference page for details about using single-user mode.

### 24.1.5.1. Multixacts and Wraparound

Multixact IDs are used to support row locking by multiple transactions. Since there is only limited space in a tuple header to store lock information, that information is encoded as a "multiple transaction ID", or multixact ID for short, whenever there is more than one transaction concurrently locking a row. Information about which transaction IDs are included in any particular multixact ID is stored separately in the \texttt{pg_multixact} subdirectory, and only the multixact ID appears in the \texttt{xmax} field in the tuple header. Like transaction IDs, multixact IDs are implemented as a 32-bit counter and corresponding storage, all of which requires careful aging management, storage cleanup, and wraparound handling. There is a separate storage area which holds the list of members in each multixact, which also uses a 32-bit counter and which must also be managed.

Whenever \texttt{VACUUM} scans any part of a table, it will replace any multixact ID it encounters which is older than \texttt{vacuum_multixact_freeze_min_age} by a different value, which can be the zero value, a single transaction ID, or a newer multixact ID. For each table, \texttt{pg_class.relminmxid} stores the oldest possible multixact ID still appearing in any tuple of that table. If this value is older than \texttt{vacuum_multixact_freeze_table_age}, an aggressive vacuum is forced. As discussed in the previous section, an aggressive vacuum means that only those pages which are known to be all-frozen will be skipped. \texttt{mxid_age()} can be used on \texttt{pg_class.relminmxid} to find its age.

Aggressive \texttt{VACUUM} scans, regardless of what causes them, enable advancing the value for that table. Eventually, as all tables in all databases are scanned and their oldest multixact values are advanced, on-disk storage for older multixacts can be removed.

As a safety device, an aggressive vacuum scan will occur for any table whose multixact-age is greater than \texttt{autovacuum_multixact_freeze_max_age}. Aggressive vacuum scans will also occur progressively for all tables, starting with those that have the oldest multixact-age, if the amount of used member storage space exceeds the amount 50% of the addressable storage space. Both of these kinds of aggressive scans will occur even if \texttt{autovacuum} is nominally disabled.

### 24.1.6. The Autovacuum Daemon

PostgreSQL has an optional but highly recommended feature called \texttt{autovacuum}, whose purpose is to automate the execution of \texttt{VACUUM} and \texttt{ANALYZE} commands. When enabled, autovacuum checks for tables that have had a large number of inserted, updated or deleted tuples. These checks use the statistics collection facility; therefore, autovacuum cannot be used unless \texttt{track_counts} is set to \texttt{true}. In the default configuration, autovacuuming is enabled and the related configuration parameters are appropriately set.
Chapter 24. Routine Database Maintenance Tasks

The “autovacuum daemon” actually consists of multiple processes. There is a persistent daemon process, called the autovacuum launcher, which is in charge of starting autovacuum worker processes for all databases. The launcher will distribute the work across time, attempting to start one worker within each database every autovacuum_naptime seconds. (Therefore, if the installation has $N$ databases, a new worker will be launched every $\text{autovacuum_naptime}/N$ seconds.) A maximum of autovacuum_max_workers worker processes are allowed to run at the same time. If there are more than autovacuum_max_workers databases to be processed, the next database will be processed as soon as the first worker finishes. Each worker process will check each table within its database and execute VACUUM and/or ANALYZE as needed. log_autovacuum_min_duration can be set to monitor autovacuum workers’ activity.

If several large tables all become eligible for vacuuming in a short amount of time, all autovacuum workers might become occupied with vacuuming those tables for a long period. This would result in other tables and databases not being vacuumed until a worker becomes available. There is no limit on how many workers might be in a single database, but workers do try to avoid repeating work that has already been done by other workers. Note that the number of running workers does not count towards max_connections or superuser_reserved_connections limits.

Tables whose relfrozenxid value is more than autovacuum_freeze_max_age transactions old are always vacuumed (this also applies to those tables whose freeze max age has been modified via storage parameters; see below). Otherwise, if the number of tuples obsoleted since the last VACUUM exceeds the “vacuum threshold”, the table is vacuumed. The vacuum threshold is defined as:

$$\text{vacuum threshold} = \text{vacuum base threshold} + \text{vacuum scale factor} \times \text{number of tuples}$$

where the vacuum base threshold is autovacuum_vacuum_threshold, the vacuum scale factor is autovacuum_vacuum_scale_factor, and the number of tuples is pg_class.reltuples. The number of obsolete tuples is obtained from the statistics collector; it is a semi-accurate count updated by each UPDATE and DELETE operation. (It is only semi-accurate because some information might be lost under heavy load.) If the relfrozenxid value of the table is more than vacuum_freeze_table_age transactions old, an aggressive vacuum is performed to freeze old tuples and advance relfrozenxid; otherwise, only pages that have been modified since the last vacuum are scanned.

For analyze, a similar condition is used: the threshold, defined as:

$$\text{analyze threshold} = \text{analyze base threshold} + \text{analyze scale factor} \times \text{number of tuples}$$

is compared to the total number of tuples inserted, updated, or deleted since the last ANALYZE.

Temporary tables cannot be accessed by autovacuum. Therefore, appropriate vacuum and analyze operations should be performed via session SQL commands.

The default thresholds and scale factors are taken from postgresql.conf, but it is possible to override them (and many other autovacuum control parameters) on a per-table basis; see Storage Parameters for more information. If a setting has been changed via a table’s storage parameters, that value is used when processing that table; otherwise the global settings are used. See Section 19.10 for more details on the global settings.

When multiple workers are running, the autovacuum cost delay parameters (see Section 19.4.4) are “balanced” among all the running workers, so that the total I/O impact on the system is the same regardless of the number of workers actually running. However, any workers processing tables whose per-table autovacuum_vacuum_cost_delay or autovacuum_vacuum_cost_limit storage parameters have been set are not considered in the balancing algorithm.

Autovacuum workers generally don’t block other commands. If a process attempts to acquire a lock that conflicts with the SHARE UPDATE EXCLUSIVE lock held by autovacuum, lock acquisition will interrupt the autovacuum. For conflicting lock modes, see Table 13-2. However, if the auto-
vacuum is running to prevent transaction ID wraparound (i.e., the autovacuum query name in the pg_stat_activity view ends with (to prevent wraparound) ), the autovacuum is not automatically interrupted.

**Warning**
Regularly running commands that acquire locks conflicting with a SHARE UPDATE EXCLUSIVE lock (e.g., ANALYZE) can effectively prevent autovacuums from ever completing.

### 24.2. Routine Reindexing

In some situations it is worthwhile to rebuild indexes periodically with the REINDEX command or a series of individual rebuilding steps.

B-tree index pages that have become completely empty are reclaimed for re-use. However, there is still a possibility of inefficient use of space: if all but a few index keys on a page have been deleted, the page remains allocated. Therefore, a usage pattern in which most, but not all, keys in each range are eventually deleted will see poor use of space. For such usage patterns, periodic reindexing is recommended.

The potential for bloat in non-B-tree indexes has not been well researched. It is a good idea to periodically monitor the index’s physical size when using any non-B-tree index type.

Also, for B-tree indexes, a freshly-constructed index is slightly faster to access than one that has been updated many times because logically adjacent pages are usually also physically adjacent in a newly built index. (This consideration does not apply to non-B-tree indexes.) It might be worthwhile to reindex periodically just to improve access speed.

REINDEX can be used safely and easily in all cases. But since the command requires an exclusive table lock, it is often preferable to execute an index rebuild with a sequence of creation and replacement steps. Index types that support CREATE INDEX with the CONCURRENTLY option can instead be recreated that way. If that is successful and the resulting index is valid, the original index can then be replaced by the newly built one using a combination of ALTER INDEX and DROP INDEX. When an index is used to enforce uniqueness or other constraints, ALTER TABLE might be necessary to swap the existing constraint with one enforced by the new index. Review this alternate multistep rebuild approach carefully before using it as there are limitations on which indexes can be reindexed this way, and errors must be handled.

### 24.3. Log File Maintenance

It is a good idea to save the database server’s log output somewhere, rather than just discarding it via /dev/null. The log output is invaluable when diagnosing problems. However, the log output tends to be voluminous (especially at higher debug levels) so you won’t want to save it indefinitely. You need to rotate the log files so that new log files are started and old ones removed after a reasonable period of time.

If you simply direct the stderr of postgres into a file, you will have log output, but the only way to truncate the log file is to stop and restart the server. This might be acceptable if you are using
PostgreSQL in a development environment, but few production servers would find this behavior acceptable.

A better approach is to send the server’s stderr output to some type of log rotation program. There is a built-in log rotation facility, which you can use by setting the configuration parameter `logging_collector` to `true` in `postgresql.conf`. The control parameters for this program are described in Section 19.8.1. You can also use this approach to capture the log data in machine readable CSV (comma-separated values) format.

Alternatively, you might prefer to use an external log rotation program if you have one that you are already using with other server software. For example, the rotatelogs tool included in the Apache distribution can be used with PostgreSQL. To do this, just pipe the server’s stderr output to the desired program. If you start the server with `pg_ctl`, then `stderr` is already redirected to stdout, so you just need a pipe command, for example:

```
pg_ctl start | rotatelogs /var/log/pgsql_log 86400
```

Another production-grade approach to managing log output is to send it to syslog and let syslog deal with file rotation. To do this, set the configuration parameter `log_destination` to `syslog` (to log to syslog only) in `postgresql.conf`. Then you can send a `SIGHUP` signal to the syslog daemon whenever you want to force it to start writing a new log file. If you want to automate log rotation, the logrotate program can be configured to work with log files from syslog.

On many systems, however, syslog is not very reliable, particularly with large log messages; it might truncate or drop messages just when you need them the most. Also, on Linux, syslog will flush each message to disk, yielding poor performance. (You can use a “-” at the start of the file name in the syslog configuration file to disable syncing.)

Note that all the solutions described above take care of starting new log files at configurable intervals, but they do not handle deletion of old, no-longer-useful log files. You will probably want to set up a batch job to periodically delete old log files. Another possibility is to configure the rotation program so that old log files are overwritten cyclically.

pgBadger\(^2\) is an external project that does sophisticated log file analysis. check_postgres\(^3\) provides Nagios alerts when important messages appear in the log files, as well as detection of many other extraordinary conditions.

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2. [https://pgbadger.darold.net/](https://pgbadger.darold.net/)
3. [https://bucardo.org/check_postgres/](https://bucardo.org/check_postgres/)
Chapter 25. Backup and Restore

As with everything that contains valuable data, PostgreSQL databases should be backed up regularly. While the procedure is essentially simple, it is important to have a clear understanding of the underlying techniques and assumptions.

There are three fundamentally different approaches to backing up PostgreSQL data:

- SQL dump
- File system level backup
- Continuous archiving

Each has its own strengths and weaknesses; each is discussed in turn in the following sections.

25.1. SQL Dump

The idea behind this dump method is to generate a file with SQL commands that, when fed back to the server, will recreate the database in the same state as it was at the time of the dump. PostgreSQL provides the utility program pg_dump for this purpose. The basic usage of this command is:

```
pg_dump dbname > dumpfile
```

As you see, pg_dump writes its result to the standard output. We will see below how this can be useful. While the above command creates a text file, pg_dump can create files in other formats that allow for parallelism and more fine-grained control of object restoration.

pg_dump is a regular PostgreSQL client application (albeit a particularly clever one). This means that you can perform this backup procedure from any remote host that has access to the database. But remember that pg_dump does not operate with special permissions. In particular, it must have read access to all tables that you want to back up, so in order to back up the entire database you almost always have to run it as a database superuser. (If you do not have sufficient privileges to back up the entire database, you can still back up portions of the database to which you do have access using options such as `-n` `schema` or `-t` `table`.)

To specify which database server pg_dump should contact, use the command line options `-h` `host` and `-p` `port`. The default host is the local host or whatever your `PGHOST` environment variable specifies. Similarly, the default port is indicated by the `PGPORT` environment variable or, failing that, by the compiled-in default. (Conveniently, the server will normally have the same compiled-in default.)

Like any other PostgreSQL client application, pg_dump will by default connect with the database user name that is equal to the current operating system user name. To override this, either specify the `-U` option or set the environment variable `PGUSER`. Remember that pg_dump connections are subject to the normal client authentication mechanisms (which are described in Chapter 20).

An important advantage of pg_dump over the other backup methods described later is that pg_dump’s output can generally be re-loaded into newer versions of PostgreSQL, whereas file-level backups and continuous archiving are both extremely server-version-specific. pg_dump is also the only method that will work when transferring a database to a different machine architecture, such as going from a 32-bit to a 64-bit server.

Dumps created by pg_dump are internally consistent, meaning, the dump represents a snapshot of the database at the time pg_dump began running. pg_dump does not block other operations on the database while it is working. (Exceptions are those operations that need to operate with an exclusive lock, such as most forms of `ALTER TABLE`.)
25.1.1. Restoring the Dump

Text files created by pg_dump are intended to be read in by the psql program. The general command form to restore a dump is

```
psql dbname < dumpfile
```

where `dumpfile` is the file output by the pg_dump command. The database `dbname` will not be created by this command, so you must create it yourself from `template0` before executing `psql` (e.g., with `createdb -T template0 dbname`). `psql` supports options similar to `pg_dump` for specifying the database server to connect to and the user name to use. See the psql reference page for more information. Non-text file dumps are restored using the `pg_restore` utility.

Before restoring an SQL dump, all the users who own objects or were granted permissions on objects in the dumped database must already exist. If they do not, the restore will fail to recreate the objects with the original ownership and/or permissions. (Sometimes this is what you want, but usually it is not.)

By default, the `psql` script will continue to execute after an SQL error is encountered. You might wish to run `psql` with the `ON_ERROR_STOP` variable set to alter that behavior and have `psql` exit with an exit status of 3 if an SQL error occurs:

```
psql --set ON_ERROR_STOP=on dbname < dumpfile
```

Either way, you will only have a partially restored database. Alternatively, you can specify that the whole dump should be restored as a single transaction, so the restore is either fully completed or fully rolled back. This mode can be specified by passing the `-1` or `--single-transaction` command-line options to `psql`. When using this mode, be aware that even a minor error can rollback a restore that has already run for many hours. However, that might still be preferable to manually cleaning up a complex database after a partially restored dump.

The ability of `pg_dump` and `psql` to write to or read from pipes makes it possible to dump a database directly from one server to another, for example:

```
pg_dump -h host1 dbname | psql -h host2 dbname
```

**Important:** The dumps produced by `pg_dump` are relative to `template0`. This means that any languages, procedures, etc. added via `template1` will also be dumped by `pg_dump`. As a result, when restoring, if you are using a customized `template1`, you must create the empty database from `template0`, as in the example above.

After restoring a backup, it is wise to run `ANALYZE` on each database so the query optimizer has useful statistics; see Section 24.1.3 and Section 24.1.6 for more information. For more advice on how to load large amounts of data into PostgreSQL efficiently, refer to Section 14.4.

25.1.2. Using pg_dumpall

`pg_dump` dumps only a single database at a time, and it does not dump information about roles or tablespaces (because those are cluster-wide rather than per-database). To support convenient dumping of the entire contents of a database cluster, the `pg_dumpall` program is provided. `pg_dumpall` backs
up each database in a given cluster, and also preserves cluster-wide data such as role and tablespace definitions. The basic usage of this command is:

```
p_g_dumpall > dumpfile
```

The resulting dump can be restored with `psql`:

```
psql -f dumpfile postgres
```

(Actually, you can specify any existing database name to start from, but if you are loading into an empty cluster then `postgres` should usually be used.) It is always necessary to have database superuser access when restoring a `pg_dumpall` dump, as that is required to restore the role and tablespace information. If you use tablespaces, make sure that the tablespace paths in the dump are appropriate for the new installation.

`pg_dumpall` works by emitting commands to re-create roles, tablespaces, and empty databases, then invoking `pg_dump` for each database. This means that while each database will be internally consistent, the snapshots of different databases are not synchronized.

Cluster-wide data can be dumped alone using the `pg_dumpall --globals-only` option. This is necessary to fully backup the cluster if running the `pg_dump` command on individual databases.

### 25.1.3. Handling Large Databases

Some operating systems have maximum file size limits that cause problems when creating large `pg_dump` output files. Fortunately, `pg_dump` can write to the standard output, so you can use standard Unix tools to work around this potential problem. There are several possible methods:

**Use compressed dumps.** You can use your favorite compression program, for example `gzip`:

```
p_g_dump dbname | gzip > filename.gz
```

Reload with:

```
gunzip -c filename.gz | psql dbname
```

or:

```
cat filename.gz | gunzip | psql dbname
```

**Use `split`.** The `split` command allows you to split the output into smaller files that are acceptable in size to the underlying file system. For example, to make chunks of 1 megabyte:

```
p_g_dump dbname | split -b 1m - filename
```

Reload with:

```
cat filename* | psql dbname
```

**Use `pg_dump`'s custom dump format.** If PostgreSQL was built on a system with the zlib compression library installed, the custom dump format will compress data as it writes it to the output file. This will produce dump file sizes similar to using `gzip`, but it has the added advantage that tables can be restored selectively. The following command dumps a database using the custom dump format:

```
p_g_dump -Fc dbname > filename
```
A custom-format dump is not a script for psql, but instead must be restored with pg_restore, for example:

```
pg_restore -d dbname filename
```

See the pg_dump and pg_restore reference pages for details.

For very large databases, you might need to combine split with one of the other two approaches.

Use **pg_dump's parallel dump feature**. To speed up the dump of a large database, you can use pg_dump's parallel mode. This will dump multiple tables at the same time. You can control the degree of parallelism with the `-j` parameter. Parallel dumps are only supported for the "directory" archive format.

```
pg_dump -j num -F d -f out.dir dbname
```

You can use `pg_restore -j` to restore a dump in parallel. This will work for any archive of either the "custom" or the "directory" archive mode, whether or not it has been created with `pg_dump -j`.

### 25.2. File System Level Backup

An alternative backup strategy is to directly copy the files that PostgreSQL uses to store the data in the database; Section 18.2 explains where these files are located. You can use whatever method you prefer for doing file system backups; for example:

```
tar -cf backup.tar /usr/local/pgsql/data
```

There are two restrictions, however, which make this method impractical, or at least inferior to the pg_dump method:

1. The database server **must** be shut down in order to get a usable backup. Half-way measures such as disallowing all connections will **not** work (in part because tar and similar tools do not take an atomic snapshot of the state of the file system, but also because of internal buffering within the server). Information about stopping the server can be found in Section 18.5. Needless to say, you also need to shut down the server before restoring the data.

2. If you have dug into the details of the file system layout of the database, you might be tempted to try to back up or restore only certain individual tables or databases from their respective files or directories. This will **not** work because the information contained in these files is not usable without the commit log files, `pg_clog/*`, which contain the commit status of all transactions. A table file is only usable with this information. Of course it is also impossible to restore only a table and the associated `pg_clog` data because that would render all other tables in the database cluster useless. So file system backups only work for complete backup and restoration of an entire database cluster.

An alternative file-system backup approach is to make a "consistent snapshot" of the data directory, if the file system supports that functionality (and you are willing to trust that it is implemented correctly). The typical procedure is to make a "frozen snapshot" of the volume containing the database, then copy the whole data directory (not just parts, see above) from the snapshot to a backup device, then release the frozen snapshot. This will work even while the database server is running. However, a backup
created in this way saves the database files in a state as if the database server was not properly shut down; therefore, when you start the database server on the backed-up data, it will think the previous server instance crashed and will replay the WAL log. This is not a problem; just be aware of it (and be sure to include the WAL files in your backup). You can perform a **CHECKPOINT** before taking the snapshot to reduce recovery time.

If your database is spread across multiple file systems, there might not be any way to obtain exactly-simultaneous frozen snapshots of all the volumes. For example, if your data files and WAL log are on different disks, or if tablespaces are on different file systems, it might not be possible to use snapshot backup because the snapshots **must** be simultaneous. Read your file system documentation very carefully before trusting the consistent-snapshot technique in such situations.

If simultaneous snapshots are not possible, one option is to shut down the database server long enough to establish all the frozen snapshots. Another option is to perform a continuous archiving base backup (Section 25.3.2) because such backups are immune to file system changes during the backup. This requires enabling continuous archiving just during the backup process; restore is done using continuous archive recovery (Section 25.3.4).

Another option is to use rsync to perform a file system backup. This is done by first running rsync while the database server is running, then shutting down the database server long enough to do an `rsync --checksum`. (--checksum is necessary because rsync only has file modification-time granularity of one second.) The second rsync will be quicker than the first, because it has relatively little data to transfer, and the end result will be consistent because the server was down. This method allows a file system backup to be performed with minimal downtime.

Note that a file system backup will typically be larger than an SQL dump. (pg_dump does not need to dump the contents of indexes for example, just the commands to recreate them.) However, taking a file system backup might be faster.

---

**25.3. Continuous Archiving and Point-in-Time Recovery (PITR)**

At all times, PostgreSQL maintains a write ahead log (WAL) in the `pg_xlog/` subdirectory of the cluster’s data directory. The log records every change made to the database’s data files. This log exists primarily for crash-safety purposes: if the system crashes, the database can be restored to consistency by “replaying” the log entries made since the last checkpoint. However, the existence of the log makes it possible to use a third strategy for backing up databases: we can combine a file-system-level backup with backup of the WAL files. If recovery is needed, we restore the file system backup and then replay from the backed-up WAL files to bring the system to a current state. This approach is more complex to administer than either of the previous approaches, but it has some significant benefits:

- We do not need a perfectly consistent file system backup as the starting point. Any internal inconsistency in the backup will be corrected by log replay (this is not significantly different from what happens during crash recovery). So we do not need a file system snapshot capability, just tar or a similar archiving tool.

- Since we can combine an indefinitely long sequence of WAL files for replay, continuous backup can be achieved simply by continuing to archive the WAL files. This is particularly valuable for large databases, where it might not be convenient to take a full backup frequently.

- It is not necessary to replay the WAL entries all the way to the end. We could stop the replay at any point and have a consistent snapshot of the database as it was at that time. Thus, this technique
Chapter 25. Backup and Restore

supports point-in-time recovery: it is possible to restore the database to its state at any time since your base backup was taken.

- If we continuously feed the series of WAL files to another machine that has been loaded with the same base backup file, we have a warm standby system: at any point we can bring up the second machine and it will have a nearly-current copy of the database.

Note: pg_dump and pg_dumpall do not produce file-system-level backups and cannot be used as part of a continuous-archiving solution. Such dumps are logical and do not contain enough information to be used by WAL replay.

As with the plain file-system-backup technique, this method can only support restoration of an entire database cluster, not a subset. Also, it requires a lot of archival storage: the base backup might be bulky, and a busy system will generate many megabytes of WAL traffic that have to be archived. Still, it is the preferred backup technique in many situations where high reliability is needed.

To recover successfully using continuous archiving (also called “online backup” by many database vendors), you need a continuous sequence of archived WAL files that extends back at least as far as the start time of your backup. So to get started, you should set up and test your procedure for archiving WAL files before you take your first base backup. Accordingly, we first discuss the mechanics of archiving WAL files.

25.3.1. Setting Up WAL Archiving

In an abstract sense, a running PostgreSQL system produces an indefinitely long sequence of WAL records. The system physically divides this sequence into WAL segment files, which are normally 16MB apiece (although the segment size can be altered when building PostgreSQL). The segment files are given numeric names that reflect their position in the abstract WAL sequence. When not using WAL archiving, the system normally creates just a few segment files and then “recycles” them by renaming no-longer-needed segment files to higher segment numbers. It’s assumed that segment files whose contents precede the checkpoint-before-last are no longer of interest and can be recycled.

When archiving WAL data, we need to capture the contents of each segment file once it is filled, and save that data somewhere before the segment file is recycled for reuse. Depending on the application and the available hardware, there could be many different ways of “saving the data somewhere”: we could copy the segment files to an NFS-mounted directory on another machine, write them onto a tape drive (ensuring that you have a way of identifying the original name of each file), or batch them together and burn them onto CDs, or something else entirely. To provide the database administrator with flexibility, PostgreSQL tries not to make any assumptions about how the archiving will be done. Instead, PostgreSQL lets the administrator specify a shell command to be executed to copy a completed segment file to wherever it needs to go. The command could be as simple as a cp, or it could invoke a complex shell script — it’s all up to you.

To enable WAL archiving, set the wal_level configuration parameter to replica or higher, archive_mode to on, and specify the shell command to use in the archive_command configuration parameter. In practice these settings will always be placed in the postgresql.conf file. In archive_command, %p is replaced by the path name of the file to archive, while %f is replaced by only the file name. (The path name is relative to the current working directory, i.e., the cluster’s data directory.) Use %% if you need to embed an actual % character in the command. The simplest useful command is something like:
Chapter 25. Backup and Restore

archive_command = 'test ! -f /mnt/server/archivedir/%f && cp %p /mnt/server/archivedir/%f'
archive_command = 'copy "%p" "C:\server\archivedir\%f"'

which will copy archivable WAL segments to the directory /mnt/server/archivedir. (This is an
example, not a recommendation, and might not work on all platforms.) After the %p and %f parameters
have been replaced, the actual command executed might look like this:

test ! -f /mnt/server/archivedir/00000001000000A900000065 && cp pg_xlog/00000001000000A9

A similar command will be generated for each new file to be archived.

The archive command will be executed under the ownership of the same user that the PostgreSQL
server is running as. Since the series of WAL files being archived contains effectively everything
in your database, you will want to be sure that the archived data is protected from prying eyes; for
example, archive into a directory that does not have group or world read access.

It is important that the archive command return zero exit status if and only if it succeeds. Upon getting
a zero result, PostgreSQL will assume that the file has been successfully archived, and will remove or
recycle it. However, a nonzero status tells PostgreSQL that the file was not archived; it will try again
periodically until it succeeds.

The archive command should generally be designed to refuse to overwrite any pre-existing archive
file. This is an important safety feature to preserve the integrity of your archive in case of administrator
error (such as sending the output of two different servers to the same archive directory).

It is advisable to test your proposed archive command to ensure that it indeed does not overwrite an
existing file, and that it returns nonzero status in this case. The example command above for Unix
ensures this by including a separate test step. On some Unix platforms, cp has switches such as -i
that can be used to do the same thing less verbosely, but you should not rely on these without verifying
that the right exit status is returned. (In particular, GNU cp will return status zero when -i is used
and the target file already exists, which is not the desired behavior.)

While designing your archiving setup, consider what will happen if the archive command fails repeat-
edly because some aspect requires operator intervention or the archive runs out of space. For example,
this could occur if you write to tape without an autochanger; when the tape fills, nothing further can
be archived until the tape is swapped. You should ensure that any error condition or request to a hu-
man operator is reported appropriately so that the situation can be resolved reasonably quickly. The
pg_xlog/ directory will continue to fill with WAL segment files until the situation is resolved. (If
the file system containing pg_xlog/ fills up, PostgreSQL will do a PANIC shutdown. No committed
transactions will be lost, but the database will remain offline until you free some space.)

The speed of the archiving command is unimportant as long as it can keep up with the average rate
at which your server generates WAL data. Normal operation continues even if the archiving process
falls a little behind. If archiving falls significantly behind, this will increase the amount of data that
would be lost in the event of a disaster. It will also mean that the pg_xlog/ directory will contain
large numbers of not-yet-archived segment files, which could eventually exceed available disk space.
You are advised to monitor the archiving process to ensure that it is working as you intend.

In writing your archive command, you should assume that the file names to be archived can be up
to 64 characters long and can contain any combination of ASCII letters, digits, and dots. It is not
necessary to preserve the original relative path (%p) but it is necessary to preserve the file name (%f).

Note that although WAL archiving will allow you to restore any modifications made to the data
in your PostgreSQL database, it will not restore changes made to configuration files (that is,
postgresql.conf, pg_hba.conf and pg_ident.conf), since those are edited manually rather
than through SQL operations. You might wish to keep the configuration files in a location that will
be backed up by your regular file system backup procedures. See Section 19.2 for how to relocate the configuration files.

The archive command is only invoked on completed WAL segments. Hence, if your server generates only little WAL traffic (or has slack periods where it does so), there could be a long delay between the completion of a transaction and its safe recording in archive storage. To put a limit on how old unarchived data can be, you can set archive_timeout to force the server to switch to a new WAL segment file at least that often. Note that archived files that are archived early due to a forced switch are still the same length as completely full files. It is therefore unwise to set a very short archive_timeout — it will bloat your archive storage. archive_timeout settings of a minute or so are usually reasonable.

Also, you can force a segment switch manually with pg_switch_xlog if you want to ensure that a just-finished transaction is archived as soon as possible. Other utility functions related to WAL management are listed in Table 9-78.

When wal_level is minimal some SQL commands are optimized to avoid WAL logging, as described in Section 14.4.7. If archiving or streaming replication were turned on during execution of one of these statements, WAL would not contain enough information for archive recovery. (Crash recovery is unaffected.) For this reason, wal_level can only be changed at server start. However, archive_command can be changed with a configuration file reload. If you wish to temporarily stop archiving, one way to do it is to set archive_command to the empty string ("""). This will cause WAL files to accumulate in pg_xlog/ until a working archive_command is re-established.

### 25.3.2. Making a Base Backup

The easiest way to perform a base backup is to use the pg_basebackup tool. It can create a base backup either as regular files or as a tar archive. If more flexibility than pg_basebackup can provide is required, you can also make a base backup using the low level API (see Section 25.3.3).

It is not necessary to be concerned about the amount of time it takes to make a base backup. However, if you normally run the server with full_page_writes disabled, you might notice a drop in performance while the backup runs since full_page_writes is effectively forced on during backup mode.

To make use of the backup, you will need to keep all the WAL segment files generated during and after the file system backup. To aid you in doing this, the base backup process creates a backup history file that is immediately stored into the WAL archive area. This file is named after the first WAL segment file that you need for the file system backup. For example, if the starting WAL file is 00000001000123400005CD the backup history file will be named something like 00000001000123400005CD.007C9330.backup. (The second part of the file name stands for an exact position within the WAL file, and can ordinarily be ignored.) Once you have safely archived the file system backup and the WAL segment files used during the backup (as specified in the backup history file), all archived WAL segments with names numerically less are no longer needed to recover the file system backup and can be deleted. However, you should consider keeping several backup sets to be absolutely certain that you can recover your data.

The backup history file is just a small text file. It contains the label string you gave to pg_basebackup, as well as the starting and ending times and WAL segments of the backup. If you used the label to identify the associated dump file, then the archived history file is enough to tell you which dump file to restore.

Since you have to keep around all the archived WAL files back to your last base backup, the interval between base backups should usually be chosen based on how much storage you want to expend on archived WAL files. You should also consider how long you are prepared to spend recovering, if
recovery should be necessary — the system will have to replay all those WAL segments, and that could take awhile if it has been a long time since the last base backup.

25.3.3. Making a Base Backup Using the Low Level API

The procedure for making a base backup using the low level APIs contains a few more steps than the `pg_basebackup` method, but is relatively simple. It is very important that these steps are executed in sequence, and that the success of a step is verified before proceeding to the next step.

Low level base backups can be made in a non-exclusive or an exclusive way. The non-exclusive method is recommended and the exclusive one is deprecated and will eventually be removed.

25.3.3.1. Making a non-exclusive low level backup

A non-exclusive low level backup is one that allows other concurrent backups to be running (both those started using the same backup API and those started using `pg_basebackup`).

1. Ensure that WAL archiving is enabled and working.

2. Connect to the server (it does not matter which database) as a user with rights to run `pg_start_backup` (superuser, or a user who has been granted EXECUTE on the function) and issue the command:

```sql
SELECT pg_start_backup('label', false, false);
```

where `label` is any string you want to use to uniquely identify this backup operation. The connection calling `pg_start_backup` must be maintained until the end of the backup, or the backup will be automatically aborted.

By default, `pg_start_backup` can take a long time to finish. This is because it performs a checkpoint, and the I/O required for the checkpoint will be spread out over a significant period of time, by default half your inter-checkpoint interval (see the configuration parameter `checkpoint_completion_target`). This is usually what you want, because it minimizes the impact on query processing. If you want to start the backup as soon as possible, change the second parameter to `true`, which will issue an immediate checkpoint using as much I/O as available.

The third parameter being `false` tells `pg_start_backup` to initiate a non-exclusive base backup.

3. Perform the backup, using any convenient file-system-backup tool such as tar or cpio (not `pg_dump` or `pg_dumpall`). It is neither necessary nor desirable to stop normal operation of the database while you do this. See Section 25.3.3.3 for things to consider during this backup.

4. In the same connection as before, issue the command:

```sql
SELECT * FROM pg_stop_backup(false);
```

This terminates backup mode. On a primary, it also performs an automatic switch to the next WAL segment. On a standby, it is not possible to automatically switch WAL segments, so you may wish to run `pg_switch_xlog` on the primary to perform a manual switch. The reason for the switch is to arrange for the last WAL segment file written during the backup interval to be ready to archive.

The `pg_stop_backup` will return one row with three values. The second of these fields should be written to a file named `backup_label` in the root directory of the backup. The third field should be written to a file named `tablespace_map` unless the field is empty. These files are vital to the backup working, and must be written without modification.
5. Once the WAL segment files active during the backup are archived, you are done. The file identified by `pg_stop_backup`’s first return value is the last segment that is required to form a complete set of backup files. On a primary, if `archive_mode` is enabled, `pg_stop_backup` does not return until the last segment has been archived. Archiving of these files happens automatically since you have already configured `archive_command`. In most cases this happens quickly, but you are advised to monitor your archive system to ensure there are no delays. If the archive process has fallen behind because of failures of the archive command, it will keep retrying until the archive succeeds and the backup is complete. If you wish to place a time limit on the execution of `pg_stop_backup`, set an appropriate `statement_timeout` value, but make note that if `pg_stop_backup` terminates because of this your backup may not be valid.

Note that on a standby `pg_stop_backup` does not wait for WAL segments to be archived so the backup process must ensure that all WAL segments required for the backup are successfully archived.

### 25.3.3.2. Making an exclusive low level backup

The process for an exclusive backup is mostly the same as for a non-exclusive one, but it differs in a few key steps. This type of backup can only be taken on a primary and does not allow concurrent backups. Prior to PostgreSQL 9.6, this was the only low-level method available, but it is now recommended that all users upgrade their scripts to use non-exclusive backups if possible.

1. Ensure that WAL archiving is enabled and working.

2. Connect to the server (it does not matter which database) as a user with rights to run `pg_start_backup` (superuser, or a user who has been granted EXECUTE on the function) and issue the command:

   ```sql
   SELECT pg_start_backup('label');
   ``

   where `label` is any string you want to use to uniquely identify this backup operation.

   `pg_start_backup` creates a `backup label` file, called `backup_label`, in the cluster directory with information about your backup, including the start time and label string. The function also creates a `tablespace map` file, called `tablespace_map`, in the cluster directory with information about tablespace symbolic links in `pg_tblspc/` if one or more such link is present. Both files are critical to the integrity of the backup, should you need to restore from it.

   By default, `pg_start_backup` can take a long time to finish. This is because it performs a checkpoint, and the I/O required for the checkpoint will be spread out over a significant period of time, by default half your inter-checkpoint interval (see the configuration parameter `checkpoint_completion_target`). This is usually what you want, because it minimizes the impact on query processing. If you want to start the backup as soon as possible, use:

   ```sql
   SELECT pg_start_backup('label', true);
   ``

   This forces the checkpoint to be done as quickly as possible.

3. Perform the backup, using any convenient file-system-backup tool such as tar or cpio (not `pg_dump` or `pg_dumpall`). It is neither necessary nor desirable to stop normal operation of the database while you do this. See Section 25.3.3.3 for things to consider during this backup.

   Note that if the server crashes during the backup it may not be possible to restart until the `backup_label` file has been manually deleted from the `PGDATA` directory.

4. Again connect to the database as a user with rights to run `pg_stop_backup` (superuser, or a user who has been granted EXECUTE on the function), and issue the command:
SELECT pg_stop_backup();

This terminates the backup mode and performs an automatic switch to the next WAL segment. The reason for the switch is to arrange for the last WAL segment file written during the backup interval to be ready to archive.

5. Once the WAL segment files active during the backup are archived, you are done. The file identified by `pg_stop_backup`’s result is the last segment that is required to form a complete set of backup files. If `archive_mode` is enabled, `pg_stop_backup` does not return until the last segment has been archived. Archiving of these files happens automatically since you have already configured `archive_command`. In most cases this happens quickly, but you are advised to monitor your archive system to ensure there are no delays. If the archive process has fallen behind because of failures of the archive command, it will keep retrying until the archive succeeds and the backup is complete. If you wish to place a time limit on the execution of `pg_stop_backup`, set an appropriate `statement_timeout` value, but make note that if `pg_stop_backup` terminates because of this your backup may not be valid.

### 25.3.3.3. Backing up the data directory

Some file system backup tools emit warnings or errors if the files they are trying to copy change while the copy proceeds. When taking a base backup of an active database, this situation is normal and not an error. However, you need to ensure that you can distinguish complaints of this sort from real errors. For example, some versions of rsync return a separate exit code for “vanished source files”, and you can write a driver script to accept this exit code as a non-error case. Also, some versions of GNU tar return an error code indistinguishable from a fatal error if a file was truncated while tar was copying it. Fortunately, GNU tar versions 1.16 and later exit with 1 if a file was changed during the backup, and 2 for other errors. With GNU tar version 1.23 and later, you can use the warning options `--warning=no-file-changed` `--warning=no-file-removed` to hide the related warning messages.

Be certain that your backup includes all of the files under the database cluster directory (e.g., `/usr/local/pgsql/data`). If you are using tablespaces that do not reside underneath this directory, be careful to include them as well (and be sure that your backup archives symbolic links as links, otherwise the restore will corrupt your tablespaces).

You should, however, omit from the backup the files within the cluster’s `pg_xlog/` subdirectory. This slight adjustment is worthwhile because it reduces the risk of mistakes when restoring. This is easy to arrange if `pg_xlog/` is a symbolic link pointing to someplace outside the cluster directory, which is a common setup anyway for performance reasons. You might also want to exclude `postmaster.pid` and `postmaster.opts`, which record information about the running postmaster, not about the postmaster which will eventually use this backup. (These files can confuse `pg_ctl`.)

It is often a good idea to also omit from the backup the files within the cluster’s `pg_replslot/` directory, so that replication slots that exist on the master do not become part of the backup. Otherwise, the subsequent use of the backup to create a standby may result in indefinite retention of WAL files on the standby, and possibly bloat on the master if hot standby feedback is enabled, because the clients that are using those replication slots will still be connecting to and updating the slots on the master, not the standby. Even if the backup is only intended for use in creating a new master, copying the replication slots isn’t expected to be particularly useful, since the contents of those slots will likely be badly out of date by the time the new master comes on line.

The backup label file includes the label string you gave to `pg_start_backup`, as well as the time at which `pg_start_backup` was run, and the name of the starting WAL file. In case of confusion it is therefore possible to look inside a backup file and determine exactly which backup session the dump
Chapter 25. Backup and Restore

file came from. The tablespace map file includes the symbolic link names as they exist in the directory pg_tblspc/ and the full path of each symbolic link. These files are not merely for your information; their presence and contents are critical to the proper operation of the system’s recovery process.

It is also possible to make a backup while the server is stopped. In this case, you obviously cannot use pg_start_backup or pg_stop_backup, and you will therefore be left to your own devices to keep track of which backup is which and how far back the associated WAL files go. It is generally better to follow the continuous archiving procedure above.

25.3.4. Recovering Using a Continuous Archive Backup

Okay, the worst has happened and you need to recover from your backup. Here is the procedure:

1. Stop the server, if it’s running.

2. If you have the space to do so, copy the whole cluster data directory and any tablespaces to a temporary location in case you need them later. Note that this precaution will require that you have enough free space on your system to hold two copies of your existing database. If you do not have enough space, you should at least save the contents of the cluster’s pg_xlog subdirectory, as it might contain logs which were not archived before the system went down.

3. Remove all existing files and subdirectories under the cluster data directory and under the root directories of any tablespaces you are using.

4. Restore the database files from your file system backup. Be sure that they are restored with the right ownership (the database system user, not root!) and with the right permissions. If you are using tablespaces, you should verify that the symbolic links in pg_tblspc/ were correctly restored.

5. Remove any files present in pg_xlog/; these came from the file system backup and are therefore probably obsolete rather than current. If you didn’t archive pg_xlog/ at all, then recreate it with proper permissions, being careful to ensure that you re-establish it as a symbolic link if you had it set up that way before.

6. If you have unarchived WAL segment files that you saved in step 2, copy them into pg_xlog/.
   (It is best to copy them, not move them, so you still have the unmodified files if a problem occurs and you have to start over.)

7. Create a recovery command file recovery.conf in the cluster data directory (see Chapter 27). You might also want to temporarily modify pg_hba.conf to prevent ordinary users from connecting until you are sure the recovery was successful.

8. Start the server. The server will go into recovery mode and proceed to read through the archived WAL files it needs. Should the recovery be terminated because of an external error, the server can simply be restarted and it will continue recovery. Upon completion of the recovery process, the server will rename recovery.conf to recovery.done (to prevent accidentally re-entering recovery mode later) and then commence normal database operations.

9. Inspect the contents of the database to ensure you have recovered to the desired state. If not, return to step 1. If all is well, allow your users to connect by restoring pg_hba.conf to normal.

The key part of all this is to set up a recovery configuration file that describes how you want to recover and how far the recovery should run. You can use recovery.conf.sample (normally located in the installation’s share/ directory) as a prototype. The one thing that you absolutely must specify in
recovery.conf is the restore_command, which tells PostgreSQL how to retrieve archived WAL file segments. Like the archive_command, this is a shell command string. It can contain %f, which is replaced by the name of the desired log file, and %p, which is replaced by the path name to copy the log file to. (The path name is relative to the current working directory, i.e., the cluster’s data directory.) Write % if you need to embed an actual % character in the command. The simplest useful command is something like:

```
restore_command = 'cp /mnt/server/archivedir/%f %p'
```

which will copy previously archived WAL segments from the directory /mnt/server/archivedir. Of course, you can use something much more complicated, perhaps even a shell script that requests the operator to mount an appropriate tape.

It is important that the command return nonzero exit status on failure. The command will be called requesting files that are not present in the archive; it must return nonzero when so asked. This is not an error condition. An exception is that if the command was terminated by a signal (other than SIGTERM, which is used as part of a database server shutdown) or an error by the shell (such as command not found), then recovery will abort and the server will not start up.

Not all of the requested files will be WAL segment files; you should also expect requests for files with a suffix of .history. Also be aware that the base name of the %p path will be different from %f; do not expect them to be interchangeable.

WAL segments that cannot be found in the archive will be sought in pg_xlog/; this allows use of recent un-archived segments. However, segments that are available from the archive will be used in preference to files in pg_xlog/.

Normally, recovery will proceed through all available WAL segments, thereby restoring the database to the current point in time (or as close as possible given the available WAL segments). Therefore, a normal recovery will end with a “file not found” message, the exact text of the error message depending upon your choice of restore_command. You may also see an error message at the start of recovery for a file named something like 00000001.history. This is also normal and does not indicate a problem in simple recovery situations; see Section 25.3.5 for discussion.

If you want to recover to some previous point in time (say, right before the junior DBA dropped your main transaction table), just specify the required stopping point in recovery.conf. You can specify the stop point, known as the “recovery target”, either by date/time, named restore point or by completion of a specific transaction ID. As of this writing only the date/time and named restore point options are very usable, since there are no tools to help you identify with any accuracy which transaction ID to use.

**Note:** The stop point must be after the ending time of the base backup, i.e., the end time of `pg_stop_backup`. You cannot use a base backup to recover to a time when that backup was in progress. (To recover to such a time, you must go back to your previous base backup and roll forward from there.)

If recovery finds corrupted WAL data, recovery will halt at that point and the server will not start. In such a case the recovery process could be re-run from the beginning, specifying a “recovery target” before the point of corruption so that recovery can complete normally. If recovery fails for an external reason, such as a system crash or if the WAL archive has become inaccessible, then the recovery can simply be restarted and it will restart almost from where it failed. Recovery restart works much like checkpointing in normal operation: the server periodically forces all its state to disk, and then updates the pg_control file to indicate that the already-processed WAL data need not be scanned again.
25.3.5. Timelines

The ability to restore the database to a previous point in time creates some complexities that are akin to science-fiction stories about time travel and parallel universes. For example, in the original history of the database, suppose you dropped a critical table at 5:15PM on Tuesday evening, but didn’t realize your mistake until Wednesday noon. Unfazed, you get out your backup, restore to the point-in-time 5:14PM Tuesday evening, and are up and running. In this history of the database universe, you never dropped the table. But suppose you later realize this wasn’t such a great idea, and would like to return to sometime Wednesday morning in the original history. You won’t be able to if, while your database was up-and-running, it overwrote some of the WAL segment files that led up to the time you now wish you could get back to. Thus, to avoid this, you need to distinguish the series of WAL records generated after you’ve done a point-in-time recovery from those that were generated in the original database history.

To deal with this problem, PostgreSQL has a notion of timelines. Whenever an archive recovery completes, a new timeline is created to identify the series of WAL records generated after that recovery. The timeline ID number is part of WAL segment file names so a new timeline does not overwrite the WAL data generated by previous timelines. It is in fact possible to archive many different timelines. While that might seem like a useless feature, it’s often a lifesaver. Consider the situation where you aren’t quite sure what point-in-time to recover to, and so have to do several point-in-time recoveries by trial and error until you find the best place to branch off from the old history. Without timelines this process would soon generate an unmanageable mess. With timelines, you can recover to any prior state, including states in timeline branches that you abandoned earlier.

Every time a new timeline is created, PostgreSQL creates a “timeline history” file that shows which timeline it branched off from and when. These history files are necessary to allow the system to pick the right WAL segment files when recovering from an archive that contains multiple timelines. Therefore, they are archived into the WAL archive area just like WAL segment files. The history files are just small text files, so it’s cheap and appropriate to keep them around indefinitely (unlike the segment files which are large). You can, if you like, add comments to a history file to record your own notes about how and why this particular timeline was created. Such comments will be especially valuable when you have a thicket of different timelines as a result of experimentation.

The default behavior of recovery is to recover along the same timeline that was current when the base backup was taken. If you wish to recover into some child timeline (that is, you want to return to some state that was itself generated after a recovery attempt), you need to specify the target timeline ID in recovery.conf. You cannot recover into timelines that branched off earlier than the base backup.

25.3.6. Tips and Examples

Some tips for configuring continuous archiving are given here.

25.3.6.1. Standalone Hot Backups

It is possible to use PostgreSQL’s backup facilities to produce standalone hot backups. These are backups that cannot be used for point-in-time recovery, yet are typically much faster to backup and restore than pg_dump dumps. (They are also much larger than pg_dump dumps, so in some cases the speed advantage might be negated.)

As with base backups, the easiest way to produce a standalone hot backup is to use the pg_basebackup tool. If you include the -x parameter when calling it, all the transaction log required to use the backup will be included in the backup automatically, and no special action is required to restore the backup.
If more flexibility in copying the backup files is needed, a lower level process can be used for standalone hot backups as well. To prepare for low level standalone hot backups, set `wal_level` to `replica` or higher, `archive_mode` to `on`, and set up an `archive_command` that performs archiving only when a switch file exists. For example:

```
archive_command = 'test ! -f /var/lib/pgsql/backup_in_progress || (test ! -f /var/lib/pgsql/archive/%f && cp %p /var/lib/pgsql/archive/%f)'
```

This command will perform archiving when `/var/lib/pgsql/backup_in_progress` exists, and otherwise silently return zero exit status (allowing PostgreSQL to recycle the unwanted WAL file).

With this preparation, a backup can be taken using a script like the following:

```
touch /var/lib/pgsql/backup_in_progress
touch /var/lib/pgsql/backup_in_progress
psql -c "select pg_start_backup('hot_backup');"
tar -cf /var/lib/pgsql/backup.tar /var/lib/pgsql/data/
rm /var/lib/pgsql/backup_in_progress
tar -rf /var/lib/pgsql/backup.tar /var/lib/pgsql/archive/
```

The switch file `/var/lib/pgsql/backup_in_progress` is created first, enabling archiving of completed WAL files to occur. After the backup the switch file is removed. Archived WAL files are then added to the backup so that both base backup and all required WAL files are part of the same tar file. Please remember to add error handling to your backup scripts.

### 25.3.6.2. Compressed Archive Logs

If archive storage size is a concern, you can use `gzip` to compress the archive files:

```
archive_command = 'gzip < %p > /var/lib/pgsql/archive/%f'
```

You will then need to use `gunzip` during recovery:

```
restore_command = 'gunzip < /mnt/server/archivedir/%f > %p'
```

### 25.3.6.3. archive_command Scripts

Many people choose to use scripts to define their `archive_command`, so that their `postgresql.conf` entry looks very simple:

```
archive_command = '/local_backup_script.sh "%p" "%f"'
```

Using a separate script file is advisable any time you want to use more than a single command in the archiving process. This allows all complexity to be managed within the script, which can be written in a popular scripting language such as bash or perl.

Examples of requirements that might be solved within a script include:

- Copying data to secure off-site data storage
- Batching WAL files so that they are transferred every three hours, rather than one at a time
- Interfacing with other backup and recovery software
- Interfacing with monitoring software to report errors
Tip: When using an archive_command script, it's desirable to enable logging_collector. Any messages written to stderr from the script will then appear in the database server log, allowing complex configurations to be diagnosed easily if they fail.

25.3.7. Caveats

At this writing, there are several limitations of the continuous archiving technique. These will probably be fixed in future releases:

- Operations on hash indexes are not presently WAL-logged, so replay will not update these indexes. This will mean that any new inserts will be ignored by the index, updated rows will apparently disappear and deleted rows will still retain pointers. In other words, if you modify a table with a hash index on it then you will get incorrect query results on a standby server. When recovery completes it is recommended that you manually REINDEX each such index after completing a recovery operation.

- If a CREATE DATABASE command is executed while a base backup is being taken, and then the template database that the CREATE DATABASE copied is modified while the base backup is still in progress, it is possible that recovery will cause those modifications to be propagated into the created database as well. This is of course undesirable. To avoid this risk, it is best not to modify any template databases while taking a base backup.

- CREATE TABLESPACE commands are WAL-logged with the literal absolute path, and will therefore be replayed as tablespace creations with the same absolute path. This might be undesirable if the log is being replayed on a different machine. It can be dangerous even if the log is being replayed on the same machine, but into a new data directory: the replay will still overwrite the contents of the original tablespace. To avoid potential gotchas of this sort, the best practice is to take a new base backup after creating or dropping tablespaces.

It should also be noted that the default WAL format is fairly bulky since it includes many disk page snapshots. These page snapshots are designed to support crash recovery, since we might need to fix partially-written disk pages. Depending on your system hardware and software, the risk of partial writes might be small enough to ignore, in which case you can significantly reduce the total volume of archived logs by turning off page snapshots using the full_page_writes parameter. (Read the notes and warnings in Chapter 30 before you do so.) Turning off page snapshots does not prevent use of the logs for PITR operations. An area for future development is to compress archived WAL data by removing unnecessary page copies even when full_page_writes is on. In the meantime, administrators might wish to reduce the number of page snapshots included in WAL by increasing the checkpoint interval parameters as much as feasible.
Database servers can work together to allow a second server to take over quickly if the primary server fails (high availability), or to allow several computers to serve the same data (load balancing). Ideally, database servers could work together seamlessly. Web servers serving static web pages can be combined quite easily by merely load-balancing web requests to multiple machines. In fact, read-only database servers can be combined relatively easily too. Unfortunately, most database servers have a read/write mix of requests, and read/write servers are much harder to combine. This is because though read-only data needs to be placed on each server only once, a write to any server has to be propagated to all servers so that future read requests to those servers return consistent results.

This synchronization problem is the fundamental difficulty for servers working together. Because there is no single solution that eliminates the impact of the sync problem for all use cases, there are multiple solutions. Each solution addresses this problem in a different way, and minimizes its impact for a specific workload.

Some solutions deal with synchronization by allowing only one server to modify the data. Servers that can modify data are called read/write, master or primary servers. Servers that track changes in the master are called standby or slave servers. A standby server that cannot be connected to until it is promoted to a master server is called a warm standby server, and one that can accept connections and serves read-only queries is called a hot standby server.

Some solutions are synchronous, meaning that a data-modifying transaction is not considered committed until all servers have committed the transaction. This guarantees that a failover will not lose any data and that all load-balanced servers will return consistent results no matter which server is queried.

In contrast, asynchronous solutions allow some delay between the time of a commit and its propagation to the other servers, opening the possibility that some transactions might be lost in the switch to a backup server, and that load balanced servers might return slightly stale results. Asynchronous communication is used when synchronous would be too slow.

Solutions can also be categorized by their granularity. Some solutions can deal only with an entire database server, while others allow control at the per-table or per-database level.

Performance must be considered in any choice. There is usually a trade-off between functionality and performance. For example, a fully synchronous solution over a slow network might cut performance by more than half, while an asynchronous one might have a minimal performance impact.

The remainder of this section outlines various failover, replication, and load balancing solutions.

### 26.1. Comparison of Different Solutions

#### Shared Disk Failover

Shared disk failover avoids synchronization overhead by having only one copy of the database. It uses a single disk array that is shared by multiple servers. If the main database server fails, the standby server is able to mount and start the database as though it were recovering from a database crash. This allows rapid failover with no data loss.

Shared hardware functionality is common in network storage devices. Using a network file system is also possible, though care must be taken that the file system has full POSIX behavior (see Section 18.2.2). One significant limitation of this method is that if the shared disk array fails or
becomes corrupt, the primary and standby servers are both nonfunctional. Another issue is that the standby server should never access the shared storage while the primary server is running.

File System (Block-Device) Replication

A modified version of shared hardware functionality is file system replication, where all changes to a file system are mirrored to a file system residing on another computer. The only restriction is that the mirroring must be done in a way that ensures the standby server has a consistent copy of the file system — specifically, writes to the standby must be done in the same order as those on the master. DRBD is a popular file system replication solution for Linux.

Transaction Log Shipping

Warm and hot standby servers can be kept current by reading a stream of write-ahead log (WAL) records. If the main server fails, the standby contains almost all of the data of the main server, and can be quickly made the new master database server. This can be synchronous or asynchronous and can only be done for the entire database server.

A standby server can be implemented using file-based log shipping (Section 26.2) or streaming replication (see Section 26.2.5), or a combination of both. For information on hot standby, see Section 26.5.

Trigger-Based Master-Standby Replication

A master-stanby replication setup sends all data modification queries to the master server. The master server asynchronously sends data changes to the standby server. The standby can answer read-only queries while the master server is running. The standby server is ideal for data warehouse queries.

Slony-I is an example of this type of replication, with per-table granularity, and support for multiple standby servers. Because it updates the standby server asynchronously (in batches), there is possible data loss during fail over.

Statement-Based Replication Middleware

With statement-based replication middleware, a program intercepts every SQL query and sends it to one or all servers. Each server operates independently. Read-write queries must be sent to all servers, so that every server receives any changes. But read-only queries can be sent to just one server, allowing the read workload to be distributed among them.

If queries are simply broadcast unmodified, functions like random(), CURRENT_TIMESTAMP, and sequences can have different values on different servers. This is because each server operates independently, and because SQL queries are broadcast (and not actual modified rows). If this is unacceptable, either the middleware or the application must query such values from a single server and then use those values in write queries. Another option is to use this replication option with a traditional master-standby setup, i.e. data modification queries are sent only to the master and are propagated to the standby servers via master-standby replication, not by the replication middleware. Care must also be taken that all transactions either commit or abort on all servers, perhaps using two-phase commit (PREPARE TRANSACTION and COMMIT PREPARED). Pgpool-II and Continuent Tungsten are examples of this type of replication.

Asynchronous Multimaster Replication

For servers that are not regularly connected or have slow communication links, like laptops or remote servers, keeping data consistent among servers is a challenge. Using asynchronous multimaster replication, each server works independently, and periodically communicates with the other servers to identify conflicting transactions. The conflicts can be resolved by users or conflict resolution rules. Bucardo is an example of this type of replication.
Chapter 26. High Availability, Load Balancing, and Replication

Synchronous Multimaster Replication

In synchronous multimaster replication, each server can accept write requests, and modified data is transmitted from the original server to every other server before each transaction commits. Heavy write activity can cause excessive locking and commit delays, leading to poor performance. Read requests can be sent to any server. Some implementations use shared disk to reduce the communication overhead. Synchronous multimaster replication is best for mostly read workloads, though its big advantage is that any server can accept write requests — there is no need to partition workloads between master and standby servers, and because the data changes are sent from one server to another, there is no problem with non-deterministic functions like \texttt{random()}. PostgreSQL does not offer this type of replication, though PostgreSQL two-phase commit (\texttt{PREPARE TRANSACTION} and \texttt{COMMIT PREPARED}) can be used to implement this in application code or middleware.

Commercial Solutions

Because PostgreSQL is open source and easily extended, a number of companies have taken PostgreSQL and created commercial closed-source solutions with unique failover, replication, and load balancing capabilities.

Table 26-1 summarizes the capabilities of the various solutions listed above.

Table 26-1. High Availability, Load Balancing, and Replication Feature Matrix

<table>
<thead>
<tr>
<th>Feature</th>
<th>Shared Disk Failover</th>
<th>File System Replication</th>
<th>Transaction Log Shipping</th>
<th>Trigger-Based Master-Standby Replication</th>
<th>Statement-Based Replication Middleware</th>
<th>Asynchronous Multimaster Replication</th>
<th>Synchronous Multimaster Replication</th>
</tr>
</thead>
<tbody>
<tr>
<td>Most Common Implementation</td>
<td>NAS</td>
<td>DRBD</td>
<td>Streaming Repl.</td>
<td>Slony</td>
<td>pgpool-II</td>
<td>Bucardo</td>
<td></td>
</tr>
<tr>
<td>Communication Method</td>
<td>shared disk</td>
<td>disk blocks</td>
<td>WAL</td>
<td>table rows</td>
<td>SQL</td>
<td>table rows</td>
<td>table rows and row locks</td>
</tr>
<tr>
<td>No special hardware required</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>Allows multiple master servers</td>
<td></td>
<td></td>
<td></td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>No master server overhead</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td></td>
<td></td>
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</tr>
</tbody>
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645
### Chapter 26. High Availability, Load Balancing, and Replication

<table>
<thead>
<tr>
<th>Feature</th>
<th>Shared Disk Failover</th>
<th>File System Replication</th>
<th>Transaction Log Shipping</th>
<th>Trigger-Based Master-Standby Replication</th>
<th>Statement-Based Replication Middleware</th>
<th>Asynchronous Multi-master Replication</th>
<th>Synchronous Multi-master Replication</th>
</tr>
</thead>
<tbody>
<tr>
<td>No waiting for multiple servers</td>
<td>•</td>
<td></td>
<td>with sync off</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>Master failure will never lose data</td>
<td>•</td>
<td>•</td>
<td>with sync on</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>Standby accept read-only queries</td>
<td></td>
<td></td>
<td>with hot</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>Per-table granularity</td>
<td></td>
<td></td>
<td></td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
<tr>
<td>No conflict resolution necessary</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
<td>•</td>
</tr>
</tbody>
</table>

There are a few solutions that do not fit into the above categories:

**Data Partitioning**

Data partitioning splits tables into data sets. Each set can be modified by only one server. For example, data can be partitioned by offices, e.g., London and Paris, with a server in each office. If queries combining London and Paris data are necessary, an application can query both servers, or master/standby replication can be used to keep a read-only copy of the other office’s data on each server.

**Multiple-Server Parallel Query Execution**

Many of the above solutions allow multiple servers to handle multiple queries, but none allow a single query to use multiple servers to complete faster. This solution allows multiple servers to work concurrently on a single query. It is usually accomplished by splitting the data among servers and having each server execute its part of the query and return results to a central server where they are combined and returned to the user. Pgpool-II has this capability. Also, this can be implemented using the PL/Proxy tool set.

### 26.2. Log-Shipping Standby Servers

Continuous archiving can be used to create a high availability (HA) cluster configuration with one or more standby servers ready to take over operations if the primary server fails. This capability is widely referred to as warm standby or log shipping.
Chapter 26. High Availability, Load Balancing, and Replication

The primary and standby server work together to provide this capability, though the servers are only loosely coupled. The primary server operates in continuous archiving mode, while each standby server operates in continuous recovery mode, reading the WAL files from the primary. No changes to the database tables are required to enable this capability, so it offers low administration overhead compared to some other replication solutions. This configuration also has relatively low performance impact on the primary server.

Directly moving WAL records from one database server to another is typically described as log shipping. PostgreSQL implements file-based log shipping by transferring WAL records one file (WAL segment) at a time. WAL files (16MB) can be shipped easily and cheaply over any distance, whether it be to an adjacent system, another system at the same site, or another system on the far side of the globe. The bandwidth required for this technique varies according to the transaction rate of the primary server. Record-based log shipping is more granular and streams WAL changes incrementally over a network connection (see Section 26.2.5).

It should be noted that log shipping is asynchronous, i.e., the WAL records are shipped after transaction commit. As a result, there is a window for data loss should the primary server suffer a catastrophic failure; transactions not yet shipped will be lost. The size of the data loss window in file-based log shipping can be limited by use of the `archive_timeout` parameter, which can be set as low as a few seconds. However such a low setting will substantially increase the bandwidth required for file shipping. Streaming replication (see Section 26.2.5) allows a much smaller window of data loss.

Recovery performance is sufficiently good that the standby will typically be only moments away from full availability once it has been activated. As a result, this is called a warm standby configuration which offers high availability. Restoring a server from an archived base backup and rollforward will take considerably longer, so that technique only offers a solution for disaster recovery, not high availability. A standby server can also be used for read-only queries, in which case it is called a Hot Standby server. See Section 26.5 for more information.

26.2.1. Planning

It is usually wise to create the primary and standby servers so that they are as similar as possible, at least from the perspective of the database server. In particular, the path names associated with tablespaces will be passed across unmodified, so both primary and standby servers must have the same mount paths for tablespaces if that feature is used. Keep in mind that if `CREATE TABLESPACE` is executed on the primary, any new mount point needed for it must be created on the primary and all standby servers before the command is executed. Hardware need not be exactly the same, but experience shows that maintaining two identical systems is easier than maintaining two dissimilar ones over the lifetime of the application and system. In any case the hardware architecture must be the same — shipping from, say, a 32-bit to a 64-bit system will not work.

In general, log shipping between servers running different major PostgreSQL release levels is not possible. It is the policy of the PostgreSQL Global Development Group not to make changes to disk formats during minor release upgrades, so it is likely that running different minor release levels on primary and standby servers will work successfully. However, no formal support for that is offered and you are advised to keep primary and standby servers at the same release level as much as possible. When updating to a new minor release, the safest policy is to update the standby servers first — a new minor release is more likely to be able to read WAL files from a previous minor release than vice versa.
26.2.2. Standby Server Operation

In standby mode, the server continuously applies WAL received from the master server. The standby server can read WAL from a WAL archive (see restore_command) or directly from the master over a TCP connection (streaming replication). The standby server will also attempt to restore any WAL found in the standby cluster’s pg_xlog directory. That typically happens after a server restart, when the standby replays again WAL that was streamed from the master before the restart, but you can also manually copy files to pg_xlog at any time to have them replayed.

At startup, the standby begins by restoring all WAL available in the archive location, calling restore_command. Once it reaches the end of WAL available there and restore_command fails, it tries to restore any WAL available in the pg_xlog directory. If that fails, and streaming replication has been configured, the standby tries to connect to the primary server and start streaming WAL from the last valid record found in archive or pg_xlog. If that fails or streaming replication is not configured, or if the connection is later disconnected, the standby goes back to step 1 and tries to restore the file from the archive again. This loop of retries from the archive, pg_xlog, and via streaming replication goes on until the server is stopped or failover is triggered by a trigger file.

Standby mode is exited and the server switches to normal operation when pg_ctl promote is run or a trigger file is found (trigger_file). Before failover, any WAL immediately available in the archive or in pg_xlog will be restored, but no attempt is made to connect to the master.

26.2.3. Preparing the Master for Standby Servers

Set up continuous archiving on the primary to an archive directory accessible from the standby, as described in Section 25.3. The archive location should be accessible from the standby even when the master is down, i.e. it should reside on the standby server itself or another trusted server, not on the master server.

If you want to use streaming replication, set up authentication on the primary server to allow replication connections from the standby server(s); that is, create a role and provide a suitable entry or entries in pg_hba.conf with the database field set to replication. Also ensure max_wal_senders is set to a sufficiently large value in the configuration file of the primary server. If replication slots will be used, ensure that max_replication_slots is set sufficiently high as well.

Take a base backup as described in Section 25.3.2 to bootstrap the standby server.

26.2.4. Setting Up a Standby Server

To set up the standby server, restore the base backup taken from primary server (see Section 25.3.4). Create a recovery command file recovery.conf in the standby’s cluster data directory, and turn on standby_mode. Set restore_command to a simple command to copy files from the WAL archive. If you plan to have multiple standby servers for high availability purposes, set recovery_target_timeline to latest, to make the standby server follow the timeline change that occurs at failover to another standby.

Note: Do not use pg_standby or similar tools with the built-in standby mode described here. restore_command should return immediately if the file does not exist; the server will retry the command again if necessary. See Section 26.4 for using tools like pg_standby.
If you want to use streaming replication, fill in `primary_conninfo` with a libpq connection string, including the host name (or IP address) and any additional details needed to connect to the primary server. If the primary needs a password for authentication, the password needs to be specified in `primary_conninfo` as well.

If you’re setting up the standby server for high availability purposes, set up WAL archiving, connections and authentication like the primary server, because the standby server will work as a primary server after failover.

If you’re using a WAL archive, its size can be minimized using the `archive_cleanup_command` parameter to remove files that are no longer required by the standby server. The `pg_archivecleanup` utility is designed specifically to be used with `archive_cleanup_command` in typical single-standby configurations, see `pg_archivecleanup`. Note however, that if you’re using the archive for backup purposes, you need to retain files needed to recover from at least the latest base backup, even if they’re no longer needed by the standby.

A simple example of a `recovery.conf` is:

```
standby_mode = 'on'
primary_conninfo = 'host=192.168.1.50 port=5432 user=foo password=foopass'
restore_command = 'cp /path/to/archive/%f %p'
archive_cleanup_command = 'pg_archivecleanup /path/to/archive %r'
```

You can have any number of standby servers, but if you use streaming replication, make sure you set `max_wal_senders` high enough in the primary to allow them to be connected simultaneously.

### 26.2.5. Streaming Replication

Streaming replication allows a standby server to stay more up-to-date than is possible with file-based log shipping. The standby connects to the primary, which streams WAL records to the standby as they’re generated, without waiting for the WAL file to be filled.

Streaming replication is asynchronous by default (see Section 26.2.8), in which case there is a small delay between committing a transaction in the primary and the changes becoming visible in the standby. This delay is however much smaller than with file-based log shipping, typically under one second assuming the standby is powerful enough to keep up with the load. With streaming replication, `archive_timeout` is not required to reduce the data loss window.

If you use streaming replication without file-based continuous archiving, the server might recycle old WAL segments before the standby has received them. If this occurs, the standby will need to be reinitialized from a new base backup. You can avoid this by setting `wal_keep_segments` to a value large enough to ensure that WAL segments are not recycled too early, or by configuring a replication slot for the standby. If you set up a WAL archive that’s accessible from the standby, these solutions are not required, since the standby can always use the archive to catch up provided it retains enough segments.

To use streaming replication, set up a file-based log-shipping standby server as described in Section 26.2. The step that turns a file-based log-shipping standby into streaming replication standby is setting `primary_conninfo` setting in the `recovery.conf` file to point to the primary server. Set `listen_addresses` and authentication options (see `pg_hba.conf`) on the primary so that the standby server can connect to the `replication` pseudo-database on the primary server (see Section 26.2.5.1).
On systems that support the keepalive socket option, setting `tcp_keepalives_idle`, `tcp_keepalives_interval` and `tcp_keepalives_count` helps the primary promptly notice a broken connection.

Set the maximum number of concurrent connections from the standby servers (see `max_wal_senders` for details).

When the standby is started and `primary_conninfo` is set correctly, the standby will connect to the primary after replaying all WAL files available in the archive. If the connection is established successfully, you will see a walreceiver process in the standby, and a corresponding walsender process in the primary.

### 26.2.5.1. Authentication

It is very important that the access privileges for replication be set up so that only trusted users can read the WAL stream, because it is easy to extract privileged information from it. Standby servers must authenticate to the primary as a superuser or an account that has the `REPLICATION` privilege. It is recommended to create a dedicated user account with `REPLICATION` and `LOGIN` privileges for replication. While `REPLICATION` privilege gives very high permissions, it does not allow the user to modify any data on the primary system, which the `SUPERUSER` privilege does.

Client authentication for replication is controlled by a `pg_hba.conf` record specifying replication in the `database` field. For example, if the standby is running on host IP `192.168.1.100` and the account name for replication is `foo`, the administrator can add the following line to the `pg_hba.conf` file on the primary:

```
# Allow the user "foo" from host 192.168.1.100 to connect to the primary
# as a replication standby if the user’s password is correctly supplied.
#
# TYPE DATABASE USER ADDRESS METHOD
host  replication foo 192.168.1.100/32 md5
```

The host name and port number of the primary, connection user name, and password are specified in the `recovery.conf` file. The password can also be set in the `~/.pgpass` file on the standby (specify replication in the `database` field). For example, if the primary is running on host IP `192.168.1.50`, port `5432`, the account name for replication is `foo`, and the password is `foopass`, the administrator can add the following line to the `recovery.conf` file on the standby:

```
# The standby connects to the primary that is running on host 192.168.1.50
# and port 5432 as the user "foo" whose password is "foopass".
primary_conninfo = 'host=192.168.1.50 port=5432 user=foo password=foopass'
```

### 26.2.5.2. Monitoring

An important health indicator of streaming replication is the amount of WAL records generated in the primary, but not yet applied in the standby. You can calculate this lag by comparing the current WAL write location on the primary with the last WAL location received by the standby. They can be retrieved using `pg_current_xlog_location` on the primary and the `pg_last_xlog_receive_location` on the standby, respectively (see Table 9-78 and Table 9-79 for details). The last WAL receive location in the standby is also displayed in the process status of the WAL receiver process, displayed using the `ps` command (see Section 28.1 for details).
You can retrieve a list of WAL sender processes via the `pg_stat_replication` view. Large differences between `pg_current_xlog_location` and `sent_location` field might indicate that the master server is under heavy load, while differences between `sent_location` and `pg_last_xlog_receive_location` on the standby might indicate network delay, or that the standby is under heavy load.

### 26.2.6. Replication Slots

Replication slots provide an automated way to ensure that the master does not remove WAL segments until they have been received by all standbys, and that the master does not remove rows which could cause a recovery conflict even when the standby is disconnected.

In lieu of using replication slots, it is possible to prevent the removal of old WAL segments using `wal_keep_segments`, or by storing the segments in an archive using `archive_command`. However, these methods often result in retaining more WAL segments than required, whereas replication slots retain only the number of segments known to be needed. An advantage of these methods is that they bound the space requirement for `pg_xlog`; there is currently no way to do this using replication slots.

Similarly, `hot_standby_feedback` and `vacuum_defer_cleanup_age` provide protection against relevant rows being removed by vacuum, but the former provides no protection during any time period when the standby is not connected, and the latter often needs to be set to a high value to provide adequate protection. Replication slots overcome these disadvantages.

#### 26.2.6.1. Querying and manipulating replication slots

Each replication slot has a name, which can contain lower-case letters, numbers, and the underscore character.

Existing replication slots and their state can be seen in the `pg_replication_slots` view.

Slots can be created and dropped either via the streaming replication protocol (see Section 51.3) or via SQL functions (see Section 9.26.6).

#### 26.2.6.2. Configuration Example

You can create a replication slot like this:

```sql
postgres=# SELECT * FROM pg_create_physical_replication_slot('node_a_slot');
   slot_name | xlog_position
------------+---------------
     node_a_slot |

postgres=# SELECT slot_name, slot_type, active FROM pg_replication_slots;
  slot_name | slot_type | active
-----------+-----------+--------
     node_a_slot | physical | f
(1 row)
```

To configure the standby to use this slot, `primary_slot_name` should be configured in the standby’s `recovery.conf`. Here is a simple example:

```bash
standby_mode = 'on'
primary_conninfo = 'host=192.168.1.50 port=5432 user=foo password=foopass'
primary_slot_name = 'node_a_slot'
```

651
26.2.7. Cascading Replication

The cascading replication feature allows a standby server to accept replication connections and stream WAL records to other standbys, acting as a relay. This can be used to reduce the number of direct connections to the master and also to minimize inter-site bandwidth overheads.

A standby acting as both a receiver and a sender is known as a cascading standby. Standbys that are more directly connected to the master are known as upstream servers, while those standby servers further away are downstream servers. Cascading replication does not place limits on the number or arrangement of downstream servers, though each standby connects to only one upstream server which eventually links to a single master/primary server.

A cascading standby sends not only WAL records received from the master but also those restored from the archive. So even if the replication connection in some upstream connection is terminated, streaming replication continues downstream for as long as new WAL records are available.

Cascading replication is currently asynchronous. Synchronous replication (see Section 26.2.8) settings have no effect on cascading replication at present.

Hot Standby feedback propagates upstream, whatever the cascaded arrangement.

If an upstream standby server is promoted to become new master, downstream servers will continue to stream from the new master if recovery_target_timeline is set to ‘latest’.

To use cascading replication, set up the cascading standby so that it can accept replication connections (that is, set max_wal_senders and hot_standby, and configure host-based authentication). You will also need to set primary_conninfo in the downstream standby to point to the cascading standby.

26.2.8. Synchronous Replication

PostgreSQL streaming replication is asynchronous by default. If the primary server crashes then some transactions that were committed may not have been replicated to the standby server, causing data loss. The amount of data loss is proportional to the replication delay at the time of failover.

Synchronous replication offers the ability to confirm that all changes made by a transaction have been transferred to one or more synchronous standby servers. This extends that standard level of durability offered by a transaction commit. This level of protection is referred to as 2-safe replication in computer science theory, and group-1-safe (group-safe and 1-safe) when synchronous_commit is set to remote_write.

When requesting synchronous replication, each commit of a write transaction will wait until confirmation is received that the commit has been written to the transaction log on disk of both the primary and standby server. The only possibility that data can be lost is if both the primary and the standby suffer crashes at the same time. This can provide a much higher level of durability, though only if the sysadmin is cautious about the placement and management of the two servers. Waiting for confirmation increases the user’s confidence that the changes will not be lost in the event of server crashes but it also necessarily increases the response time for the requesting transaction. The minimum wait time is the round-trip time between primary to standby.

Read only transactions and transaction rollbacks need not wait for replies from standby servers. Sub-transaction commits do not wait for responses from standby servers, only top-level commits. Long
running actions such as data loading or index building do not wait until the very final commit message. All two-phase commit actions require commit waits, including both prepare and commit.

### 26.2.8.1. Basic Configuration

Once streaming replication has been configured, configuring synchronous replication requires only one additional configuration step: `synchronous_standby_names` must be set to a non-empty value. `synchronous_commit` must also be set to `on`, but since this is the default value, typically no change is required. (See Section 19.5.1 and Section 19.6.2.) This configuration will cause each commit to wait for confirmation that the standby has written the commit record to durable storage. `synchronous_commit` can be set by individual users, so it can be configured in the configuration file, for particular users or databases, or dynamically by applications, in order to control the durability guarantee on a per-transaction basis.

After a commit record has been written to disk on the primary, the WAL record is then sent to the standby. The standby sends reply messages each time a new batch of WAL data is written to disk, unless `wal_receiver_status_interval` is set to zero on the standby. In the case that `synchronous_commit` is set to `remote_apply`, the standby sends reply messages when the commit record is replayed, making the transaction visible. If the standby is chosen as a synchronous standby, from a priority list of `synchronous_standby_names` on the primary, the reply messages from that standby will be considered along with those from other synchronous standbys to decide when to release transactions waiting for confirmation that the commit record has been received. These parameters allow the administrator to specify which standby servers should be synchronous standbys. Note that the configuration of synchronous replication is mainly on the master. Named standbys must be directly connected to the master; the master knows nothing about downstream standby servers using cascaded replication.

Setting `synchronous_commit` to `remote_write` will cause each commit to wait for confirmation that the standby has received the commit record and written it out to its own operating system, but not for the data to be flushed to disk on the standby. This setting provides a weaker guarantee of durability than `on` does: the standby could lose the data in the event of an operating system crash, though not a PostgreSQL crash. However, it’s a useful setting in practice because it can decrease the response time for the transaction. Data loss could only occur if both the primary and the standby crash and the database of the primary gets corrupted at the same time.

Setting `synchronous_commit` to `remote_apply` will cause each commit to wait until the current synchronous standbys report that they have replayed the transaction, making it visible to user queries. In simple cases, this allows for load balancing with causal consistency.

Users will stop waiting if a fast shutdown is requested. However, as when using asynchronous replication, the server will not fully shutdown until all outstanding WAL records are transferred to the currently connected standby servers.

### 26.2.8.2. Multiple Synchronous Standbys

Synchronous replication supports one or more synchronous standby servers; transactions will wait until all the standby servers which are considered as synchronous confirm receipt of their data. The number of synchronous standbys that transactions must wait for replies from is specified in `synchronous_standby_names`. This parameter also specifies a list of standby names, which determines the priority of each standby for being chosen as a synchronous standby. The standbys whose names appear earlier in the list are given higher priority and will be considered as synchronous. Other standby servers appearing later in this list represent potential synchronous standbys. If any of the cur-
rent synchronous standbys disconnects for whatever reason, it will be replaced immediately with the next-highest-priority standby.

An example of `synchronous_standby_names` for multiple synchronous standbys is:

```
synchronous_standby_names = '2 (s1, s2, s3)'
```

In this example, if four standby servers `s1`, `s2`, `s3` and `s4` are running, the two standbys `s1` and `s2` will be chosen as synchronous standbys because their names appear early in the list of standby names. `s3` is a potential synchronous standby and will take over the role of synchronous standby when either of `s1` or `s2` fails. `s4` is an asynchronous standby since its name is not in the list.

### 26.2.8.3. Planning for Performance

Synchronous replication usually requires carefully planned and placed standby servers to ensure applications perform acceptably. Waiting doesn’t utilize system resources, but transaction locks continue to be held until the transfer is confirmed. As a result, incautious use of synchronous replication will reduce performance for database applications because of increased response times and higher contention.

PostgreSQL allows the application developer to specify the durability level required via replication. This can be specified for the system overall, though it can also be specified for specific users or connections, or even individual transactions.

For example, an application workload might consist of: 10% of changes are important customer details, while 90% of changes are less important data that the business can more easily survive if it is lost, such as chat messages between users.

With synchronous replication options specified at the application level (on the primary) we can offer synchronous replication for the most important changes, without slowing down the bulk of the total workload. Application level options are an important and practical tool for allowing the benefits of synchronous replication for high performance applications.

You should consider that the network bandwidth must be higher than the rate of generation of WAL data.

### 26.2.8.4. Planning for High Availability

`synchronous_standby_names` specifies the number and names of synchronous standbys that transaction commits made when `synchronous_commit` is set to `on`, `remote_apply` or `remote_write` will wait for responses from. Such transaction commits may never be completed if any one of synchronous standbys should crash.

The best solution for high availability is to ensure you keep as many synchronous standbys as requested. This can be achieved by naming multiple potential synchronous standbys using `synchronous_standby_names`. The standbys whose names appear earlier in the list will be used as synchronous standbys. Standbys listed after these will take over the role of synchronous standby if one of current ones should fail.

When a standby first attaches to the primary, it will not yet be properly synchronized. This is described as `catchup` mode. Once the lag between standby and primary reaches zero for the first time we move to real-time `streaming` state. The catch-up duration may be long immediately after the standby has been created. If the standby is shut down, then the catch-up period will increase according to the length of time the standby has been down. The standby is only able to become a synchronous standby once it has reached `streaming` state.
Chapter 26. High Availability, Load Balancing, and Replication

If primary restarts while commits are waiting for acknowledgement, those waiting transactions will be marked fully committed once the primary database recovers. There is no way to be certain that all standbys have received all outstanding WAL data at time of the crash of the primary. Some transactions may not show as committed on the standby, even though they show as committed on the primary. The guarantee we offer is that the application will not receive explicit acknowledgement of the successful commit of a transaction until the WAL data is known to be safely received by all the synchronous standbys.

If you really cannot keep as many synchronous standbys as requested then you should decrease the number of synchronous standbys that transaction commits must wait for responses from in synchronous_standby_names (or disable it) and reload the configuration file on the primary server.

If the primary is isolated from remaining standby servers you should fail over to the best candidate of those other remaining standby servers.

If you need to re-create a standby server while transactions are waiting, make sure that the commands pg_start_backup() and pg_stop_backup() are run in a session with synchronous_commit = off, otherwise those requests will wait forever for the standby to appear.

26.2.9. Continuous archiving in standby

When continuous WAL archiving is used in a standby, there are two different scenarios: the WAL archive can be shared between the primary and the standby, or the standby can have its own WAL archive. When the standby has its own WAL archive, set archive_mode to always, and the standby will call the archive command for every WAL segment it receives, whether it’s by restoring from the archive or by streaming replication. The shared archive can be handled similarly, but the archive_command must test if the file being archived exists already, and if the existing file has identical contents. This requires more care in the archive_command, as it must be careful to not overwrite an existing file with different contents, but return success if the exactly same file is archived twice. And all that must be done free of race conditions, if two servers attempt to archive the same file at the same time.

If archive_mode is set to on, the archiver is not enabled during recovery or standby mode. If the standby server is promoted, it will start archiving after the promotion, but will not archive any WAL it did not generate itself. To get a complete series of WAL files in the archive, you must ensure that all WAL is archived, before it reaches the standby. This is inherently true with file-based log shipping, as the standby can only restore files that are found in the archive, but not if streaming replication is enabled. When a server is not in recovery mode, there is no difference between on and always modes.

26.3. Failover

If the primary server fails then the standby server should begin failover procedures.

If the standby server fails then no failover need take place. If the standby server can be restarted, even some time later, then the recovery process can also be restarted immediately, taking advantage of restartable recovery. If the standby server cannot be restarted, then a full new standby server instance should be created.
Chapter 26. High Availability, Load Balancing, and Replication

If the primary server fails and the standby server becomes the new primary, and then the old primary restarts, you must have a mechanism for informing the old primary that it is no longer the primary. This is sometimes known as STONITH (Shoot The Other Node In The Head), which is necessary to avoid situations where both systems think they are the primary, which will lead to confusion and ultimately data loss.

Many failover systems use just two systems, the primary and the standby, connected by some kind of heartbeat mechanism to continually verify the connectivity between the two and the viability of the primary. It is also possible to use a third system (called a witness server) to prevent some cases of inappropriate failover, but the additional complexity might not be worthwhile unless it is set up with sufficient care and rigorous testing.

PostgreSQL does not provide the system software required to identify a failure on the primary and notify the standby database server. Many such tools exist and are well integrated with the operating system facilities required for successful failover, such as IP address migration.

Once failover to the standby occurs, there is only a single server in operation. This is known as a degenerate state. The former standby is now the primary, but the former primary is down and might stay down. To return to normal operation, a standby server must be recreated, either on the former primary system when it comes up, or on a third, possibly new, system. The pg_rewind utility can be used to speed up this process on large clusters. Once complete, the primary and standby can be considered to have switched roles. Some people choose to use a third server to provide backup for the new primary until the new standby server is recreated, though clearly this complicates the system configuration and operational processes.

So, switching from primary to standby server can be fast but requires some time to re-prepare the failover cluster. Regular switching from primary to standby is useful, since it allows regular downtime on each system for maintenance. This also serves as a test of the failover mechanism to ensure that it will really work when you need it. Written administration procedures are advised.

To trigger failover of a log-shipping standby server, run 
\[
\text{pg_ctl promote}
\]
or create a trigger file with the file name and path specified by the \text{trigger_file} setting in \text{recovery.conf}. If you're planning to use \text{pg_ctl promote} to fail over, \text{trigger_file} is not required. If you're setting up the reporting servers that are only used to offload read-only queries from the primary, not for high availability purposes, you don't need to promote it.

### 26.4. Alternative Method for Log Shipping

An alternative to the built-in standby mode described in the previous sections is to use a \text{restore_command} that polls the archive location. This was the only option available in versions 8.4 and below. In this setup, set \text{standby_mode off}, because you are implementing the polling required for standby operation yourself. See the pg_standby module for a reference implementation of this.

Note that in this mode, the server will apply WAL one file at a time, so if you use the standby server for queries (see Hot Standby), there is a delay between an action in the master and when the action becomes visible in the standby, corresponding the time it takes to fill up the WAL file. \text{archive_timeout} can be used to make that delay shorter. Also note that you can’t combine streaming replication with this method.

The operations that occur on both primary and standby servers are normal continuous archiving and recovery tasks. The only point of contact between the two database servers is the archive of WAL files that both share: primary writing to the archive, standby reading from the archive. Care must be taken to ensure that WAL archives from separate primary servers do not become mixed together or confused. The archive need not be large if it is only required for standby operation.
The magic that makes the two loosely coupled servers work together is simply a `restore_command` used on the standby that, when asked for the next WAL file, waits for it to become available from the primary. The `restore_command` is specified in the `recovery.conf` file on the standby server. Normal recovery processing would request a file from the WAL archive, reporting failure if the file was unavailable. For standby processing it is normal for the next WAL file to be unavailable, so the standby must wait for it to appear. For files ending in `.history` there is no need to wait, and a non-zero return code must be returned. A waiting `restore_command` can be written as a custom script that loops after polling for the existence of the next WAL file. There must also be some way to trigger failover, which should interrupt the `restore_command`, break the loop and return a file-not-found error to the standby server. This ends recovery and the standby will then come up as a normal server.

Pseudocode for a suitable `restore_command` is:

```c
triggered = false;
while (!NextWALFileReady() && !triggered)
{
    sleep(100000L); /* wait for ~0.1 sec */
    if (CheckForExternalTrigger())
        triggered = true;
}
if (!triggered)
    CopyWALFileForRecovery();
```

A working example of a waiting `restore_command` is provided in the `pg_standby` module. It should be used as a reference on how to correctly implement the logic described above. It can also be extended as needed to support specific configurations and environments.

The method for triggering failover is an important part of planning and design. One potential option is the `restore_command` command. It is executed once for each WAL file, but the process running the `restore_command` is created and dies for each file, so there is no daemon or server process, and signals or a signal handler cannot be used. Therefore, the `restore_command` is not suitable to trigger failover. It is possible to use a simple timeout facility, especially if used in conjunction with a known `archive_timeout` setting on the primary. However, this is somewhat error prone since a network problem or busy primary server might be sufficient to initiate failover. A notification mechanism such as the explicit creation of a trigger file is ideal, if this can be arranged.

### 26.4.1. Implementation

The short procedure for configuring a standby server using this alternative method is as follows. For full details of each step, refer to previous sections as noted.

1. Set up primary and standby systems as nearly identical as possible, including two identical copies of PostgreSQL at the same release level.
2. Set up continuous archiving from the primary to a WAL archive directory on the standby server. Ensure that `archive_mode`, `archive_command` and `archive_timeout` are set appropriately on the primary (see Section 25.3.1).
3. Make a base backup of the primary server (see Section 25.3.2), and load this data onto the standby.
4. Begin recovery on the standby server from the local WAL archive, using a `recovery.conf` that specifies a `restore_command` that waits as described previously (see Section 25.3.4).
Recovery treats the WAL archive as read-only, so once a WAL file has been copied to the standby system it can be copied to tape at the same time as it is being read by the standby database server. Thus, running a standby server for high availability can be performed at the same time as files are stored for longer term disaster recovery purposes.

For testing purposes, it is possible to run both primary and standby servers on the same system. This does not provide any worthwhile improvement in server robustness, nor would it be described as HA.

### 26.4.2. Record-based Log Shipping

It is also possible to implement record-based log shipping using this alternative method, though this requires custom development, and changes will still only become visible to hot standby queries after a full WAL file has been shipped.

An external program can call the `pg_xlogfile_name_offset()` function (see Section 9.26) to find out the file name and the exact byte offset within it of the current end of WAL. It can then access the WAL file directly and copy the data from the last known end of WAL through the current end over to the standby servers. With this approach, the window for data loss is the polling cycle time of the copying program, which can be very small, and there is no wasted bandwidth from forcing partially-used segment files to be archived. Note that the standby servers’ `restore_command` scripts can only deal with whole WAL files, so the incrementally copied data is not ordinarily made available to the standby servers. It is of use only when the primary dies — then the last partial WAL file is fed to the standby before allowing it to come up. The correct implementation of this process requires cooperation of the `restore_command` script with the data copying program.

Starting with PostgreSQL version 9.0, you can use streaming replication (see Section 26.2.5) to achieve the same benefits with less effort.

### 26.5. Hot Standby

Hot Standby is the term used to describe the ability to connect to the server and run read-only queries while the server is in archive recovery or standby mode. This is useful both for replication purposes and for restoring a backup to a desired state with great precision. The term Hot Standby also refers to the ability of the server to move from recovery through to normal operation while users continue running queries and/or keep their connections open.

Running queries in hot standby mode is similar to normal query operation, though there are several usage and administrative differences explained below.

#### 26.5.1. User’s Overview

When the `hot_standby` parameter is set to true on a standby server, it will begin accepting connections once the recovery has brought the system to a consistent state. All such connections are strictly read-only; not even temporary tables may be written.

The data on the standby takes some time to arrive from the primary server so there will be a measurable delay between primary and standby. Running the same query nearly simultaneously on both primary and standby might therefore return differing results. We say that data on the standby is eventually consistent with the primary. Once the commit record for a transaction is replayed on the standby, the changes made by that transaction will be visible to any new snapshots taken on the standby. Snapshots
may be taken at the start of each query or at the start of each transaction, depending on the current transaction isolation level. For more details, see Section 13.2.

Transactions started during hot standby may issue the following commands:

- Query access - SELECT, COPY TO
- Cursor commands - DECLARE, FETCH, CLOSE
- Parameters - SHOW, SET, RESET
- Transaction management commands
  - BEGIN, END, ABORT, START TRANSACTION
  - SAVEPOINT, RELEASE, ROLLBACK TO SAVEPOINT
  - EXCEPTION blocks and other internal subtransactions
- LOCK TABLE, though only when explicitly in one of these modes: ACCESS SHARE, ROW SHARE or ROW EXCLUSIVE.
- Plans and resources - PREPARE, EXECUTE, DEALLOCATE, DISCARD
- Plugins and extensions - LOAD
- UNLISTEN

Transactions started during hot standby will never be assigned a transaction ID and cannot write to the system write-ahead log. Therefore, the following actions will produce error messages:

- Data Manipulation Language (DML) - INSERT, UPDATE, DELETE, COPY FROM, TRUNCATE. Note that there are no allowed actions that result in a trigger being executed during recovery. This restriction applies even to temporary tables, because table rows cannot be read or written without assigning a transaction ID, which is currently not possible in a Hot Standby environment.
- Data Definition Language (DDL) - CREATE, DROP, ALTER, COMMENT. This restriction applies even to temporary tables, because carrying out these operations would require updating the system catalog tables.
- SELECT ... FOR SHARE | UPDATE, because row locks cannot be taken without updating the underlying data files.
- Rules on SELECT statements that generate DML commands.
- LOCK that explicitly requests a mode higher than ROW EXCLUSIVE MODE.
- LOCK in short default form, since it requests ACCESS EXCLUSIVE MODE.
- Transaction management commands that explicitly set non-read-only state:
  - BEGIN READ WRITE, START TRANSACTION READ WRITE
  - SET TRANSACTION READ WRITE, SET SESSION CHARACTERISTICS AS TRANSACTION READ WRITE
  - SET transaction_read_only = off
- Two-phase commit commands - PREPARE TRANSACTION, COMMIT PREPARED, ROLLBACK PREPARED because even read-only transactions need to write WAL in the prepare phase (the first phase of two phase commit).
- Sequence updates - nextval(), setval()
In normal operation, “read-only” transactions are allowed to use \texttt{LISTEN} and \texttt{NOTIFY}, so Hot Standby sessions operate under slightly tighter restrictions than ordinary read-only sessions. It is possible that some of these restrictions might be loosened in a future release.

During hot standby, the parameter \texttt{transaction_read_only} is always true and may not be changed. But as long as no attempt is made to modify the database, connections during hot standby will act much like any other database connection. If failover or switchover occurs, the database will switch to normal processing mode. Sessions will remain connected while the server changes mode. Once hot standby finishes, it will be possible to initiate read-write transactions (even from a session begun during hot standby).

Users will be able to tell whether their session is read-only by issuing \texttt{SHOW transaction_read_only}. In addition, a set of functions (Table 9-79) allow users to access information about the standby server. These allow you to write programs that are aware of the current state of the database. These can be used to monitor the progress of recovery, or to allow you to write complex programs that restore the database to particular states.

### 26.5.2. Handling Query Conflicts

The primary and standby servers are in many ways loosely connected. Actions on the primary will have an effect on the standby. As a result, there is potential for negative interactions or conflicts between them. The easiest conflict to understand is performance: if a huge data load is taking place on the primary then this will generate a similar stream of WAL records on the standby, so standby queries may contend for system resources, such as I/O.

There are also additional types of conflict that can occur with Hot Standby. These conflicts are hard conflicts in the sense that queries might need to be canceled and, in some cases, sessions disconnected to resolve them. The user is provided with several ways to handle these conflicts. Conflict cases include:

- Access Exclusive locks taken on the primary server, including both explicit \texttt{LOCK} commands and various DDL actions, conflict with table accesses in standby queries.
- Dropping a tablespace on the primary conflicts with standby queries using that tablespace for temporary work files.
- Dropping a database on the primary conflicts with sessions connected to that database on the standby.
- Application of a vacuum cleanup record from WAL conflicts with standby transactions whose snapshots can still “see” any of the rows to be removed.
- Application of a vacuum cleanup record from WAL conflicts with queries accessing the target page on the standby, whether or not the data to be removed is visible.

On the primary server, these cases simply result in waiting; and the user might choose to cancel either of the conflicting actions. However, on the standby there is no choice: the WAL-logged action already occurred on the primary so the standby must not fail to apply it. Furthermore, allowing WAL application to wait indefinitely may be very undesirable, because the standby’s state will become increasingly far behind the primary’s. Therefore, a mechanism is provided to forcibly cancel standby queries that conflict with to-be-applied WAL records.
An example of the problem situation is an administrator on the primary server running `DROP TABLE` on a table that is currently being queried on the standby server. Clearly the standby query cannot continue if the `DROP TABLE` is applied on the standby. If this situation occurred on the primary, the `DROP TABLE` would wait until the other query had finished. But when `DROP TABLE` is run on the primary, the primary doesn’t have information about what queries are running on the standby, so it will not wait for any such standby queries. The WAL change records come through to the standby while the standby query is still running, causing a conflict. The standby server must either delay application of the WAL records (and everything after them, too) or else cancel the conflicting query so that the `DROP TABLE` can be applied.

When a conflicting query is short, it’s typically desirable to allow it to complete by delaying WAL application for a little bit; but a long delay in WAL application is usually not desirable. So the cancel mechanism has parameters, `max_standby_archive_delay` and `max_standby_streaming_delay`, that define the maximum allowed delay in WAL application. Conflicting queries will be canceled once it has taken longer than the relevant delay setting to apply any newly-received WAL data. There are two parameters so that different delay values can be specified for the case of reading WAL data from an archive (i.e., initial recovery from a base backup or “catching up” a standby server that has fallen far behind) versus reading WAL data via streaming replication.

In a standby server that exists primarily for high availability, it’s best to set the delay parameters relatively short, so that the server cannot fall far behind the primary due to delays caused by standby queries. However, if the standby server is meant for executing long-running queries, then a high or even infinite delay value may be preferable. Keep in mind however that a long-running query could cause other sessions on the standby server to not see recent changes on the primary, if it delays application of WAL records.

Once the delay specified by `max_standby_archive_delay` or `max_standby_streaming_delay` has been exceeded, conflicting queries will be canceled. This usually results just in a cancellation error, although in the case of replaying a `DROP DATABASE` the entire conflicting session will be terminated. Also, if the conflict is over a lock held by an idle transaction, the conflicting session is terminated (this behavior might change in the future).

Canceled queries may be retried immediately (after beginning a new transaction, of course). Since query cancellation depends on the nature of the WAL records being replayed, a query that was canceled may well succeed if it is executed again.

Keep in mind that the delay parameters are compared to the elapsed time since the WAL data was received by the standby server. Thus, the grace period allowed to any one query on the standby is never more than the delay parameter, and could be considerably less if the standby has already fallen behind as a result of waiting for previous queries to complete, or as a result of being unable to keep up with a heavy update load.

The most common reason for conflict between standby queries and WAL replay is “early cleanup”. Normally, PostgreSQL allows cleanup of old row versions when there are no transactions that need to see them to ensure correct visibility of data according to MVCC rules. However, this rule can only be applied for transactions executing on the master. So it is possible that cleanup on the master will remove row versions that are still visible to a transaction on the standby.

Experienced users should note that both row version cleanup and row version freezing will potentially conflict with standby queries. Running a manual `VACUUM FREEZE` is likely to cause conflicts even on tables with no updated or deleted rows.

Users should be clear that tables that are regularly and heavily updated on the primary server will quickly cause cancellation of longer running queries on the standby. In such cases the setting of a finite value for `max_standby_archive_delay` or `max_standby_streaming_delay` can be considered similar to setting `statement_timeout`.
Remedial possibilities exist if the number of standby-query cancellations is found to be unacceptable. The first option is to set the parameter `hot_standby_feedback`, which prevents `VACUUM` from removing recently-dead rows and so cleanup conflicts do not occur. If you do this, you should note that this will delay cleanup of dead rows on the primary, which may result in undesirable table bloat. However, the cleanup situation will be no worse than if the standby queries were running directly on the primary server, and you are still getting the benefit of off-loading execution onto the standby. If standby servers connect and disconnect frequently, you might want to make adjustments to handle the period when `hot_standby_feedback` feedback is not being provided. For example, consider increasing `max_standby Archive_delay` so that queries are not rapidly canceled by conflicts in WAL archive files during disconnected periods. You should also consider increasing `max_standby_streaming_delay` to avoid rapid cancellations by newly-arrived streaming WAL entries after reconnection.

Another option is to increase `vacuum_defer_cleanup_age` on the primary server, so that dead rows will not be cleaned up as quickly as they normally would be. This will allow more time for queries to execute before they are canceled on the standby, without having to set a high `max_standby_streaming_delay`. However it is difficult to guarantee any specific execution-time window with this approach, since `vacuum_defer_cleanup_age` is measured in transactions executed on the primary server.

The number of query cancels and the reason for them can be viewed using the `pg_stat_database_conflicts` system view on the standby server. The `pg_stat_database` system view also contains summary information.

### 26.5.3. Administrator’s Overview

If `hot_standby` is turned on in `postgresql.conf` and there is a `recovery.conf` file present, the server will run in Hot Standby mode. However, it may take some time for Hot Standby connections to be allowed, because the server will not accept connections until it has completed sufficient recovery to provide a consistent state against which queries can run. During this period, clients that attempt to connect will be refused with an error message. To confirm the server has come up, either loop trying to connect from the application, or look for these messages in the server logs:

```
LOG: entering standby mode

... then some time later ...

LOG: consistent recovery state reached
LOG: database system is ready to accept read only connections
```

Consistency information is recorded once per checkpoint on the primary. It is not possible to enable hot standby when reading WAL written during a period when `wal_level` was not set to `replica` or `logical` on the primary. Reaching a consistent state can also be delayed in the presence of both of these conditions:

- A write transaction has more than 64 subtransactions
- Very long-lived write transactions

If you are running file-based log shipping ("warm standby"), you might need to wait until the next WAL file arrives, which could be as long as the `archive_timeout` setting on the primary.
The setting of some parameters on the standby will need reconfiguration if they have been changed on the primary. For these parameters, the value on the standby must be equal to or greater than the value on the primary. If these parameters are not set high enough then the standby will refuse to start. Higher values can then be supplied and the server restarted to begin recovery again. These parameters are:

- `max_connections`
- `max_prepared_transactions`
- `max_locks_per_transaction`
- `max_worker_processes`

It is important that the administrator select appropriate settings for `max_standby_archive_delay` and `max_standby_streaming_delay`. The best choices vary depending on business priorities. For example, if the server is primarily tasked as a High Availability server, then you will want low delay settings, perhaps even zero, though that is a very aggressive setting. If the standby server is tasked as an additional server for decision support queries then it might be acceptable to set the maximum delay values to many hours, or even -1 which means wait forever for queries to complete.

Transaction status "hint bits" written on the primary are not WAL-logged, so data on the standby will likely re-write the hints again on the standby. Thus, the standby server will still perform disk writes even though all users are read-only; no changes occur to the data values themselves. Users will still write large sort temporary files and re-generate relcache info files, so no part of the database is truly read-only during hot standby mode. Note also that writes to remote databases using dblink module, and other operations outside the database using PL functions will still be possible, even though the transaction is read-only locally.

The following types of administration commands are not accepted during recovery mode:

- Data Definition Language (DDL) - e.g. `CREATE INDEX`
- Privilege and Ownership - `GRANT`, `REVOKE`, `REASSIGN`
- Maintenance commands - `ANALYZE`, `VACUUM`, `CLUSTER`, `REINDEX`

Again, note that some of these commands are actually allowed during "read only" mode transactions on the primary.

As a result, you cannot create additional indexes that exist solely on the standby, nor statistics that exist solely on the standby. If these administration commands are needed, they should be executed on the primary, and eventually those changes will propagate to the standby.

`pg_cancel_backend()` and `pg_terminate_backend()` will work on user backends, but not the Startup process, which performs recovery. `pg_stat_activity` does not show an entry for the Startup process, nor do recovering transactions show as active. As a result, `pg_prepared_xacts` is always empty during recovery. If you wish to resolve in-doubt prepared transactions, view `pg_prepared_xacts` on the primary and issue commands to resolve transactions there.

`pg_locks` will show locks held by backends, as normal. `pg_locks` also shows a virtual transaction managed by the Startup process that owns all `AccessExclusiveLocks` held by transactions being replayed by recovery. Note that the Startup process does not acquire locks to make database changes, and thus locks other than `AccessExclusiveLocks` do not show in `pg_locks` for the Startup process; they are just presumed to exist.
Chapter 26. High Availability, Load Balancing, and Replication

The Nagios plugin check_pgsql will work, because the simple information it checks for exists. The check_postgres monitoring script will also work, though some reported values could give different or confusing results. For example, last vacuum time will not be maintained, since no vacuum occurs on the standby. Vacuums running on the primary do still send their changes to the standby.

WAL file control commands will not work during recovery, e.g. pg_start_backup, pg_switch_xlog etc.

Dynamically loadable modules work, including pg_stat_statements.

Advisory locks work normally in recovery, including deadlock detection. Note that advisory locks are never WAL logged, so it is impossible for an advisory lock on either the primary or the standby to conflict with WAL replay. Nor is it possible to acquire an advisory lock on the primary and have it initiate a similar advisory lock on the standby. Advisory locks relate only to the server on which they are acquired.

Trigger-based replication systems such as Slony, Londiste and Bucardo won’t run on the standby at all, though they will run happily on the primary server as long as the changes are not sent to standby servers to be applied. WAL replay is not trigger-based so you cannot relay from the standby to any system that requires additional database writes or relies on the use of triggers.

New OIDs cannot be assigned, though some UUID generators may still work as long as they do not rely on writing new status to the database.

Currently, temporary table creation is not allowed during read only transactions, so in some cases existing scripts will not run correctly. This restriction might be relaxed in a later release. This is both a SQL Standard compliance issue and a technical issue.

DROP TABLESPACE can only succeed if the tablespace is empty. Some standby users may be actively using the tablespace via their temp_tablespaces parameter. If there are temporary files in the tablespace, all active queries are canceled to ensure that temporary files are removed, so the tablespace can be removed and WAL replay can continue.

Running DROP DATABASE or ALTER DATABASE ... SET TABLESPACE on the primary will generate a WAL entry that will cause all users connected to that database on the standby to be forcibly disconnected. This action occurs immediately, whatever the setting of max_standby_streaming_delay. Note that ALTER DATABASE ... RENAME does not disconnect users, which in most cases will go unnoticed, though might in some cases cause a program confusion if it depends in some way upon database name.

In normal (non-recovery) mode, if you issue DROP USER or DROP ROLE for a role with login capability while that user is still connected then nothing happens to the connected user - they remain connected. The user cannot reconnect however. This behavior applies in recovery also, so a DROP USER on the primary does not disconnect that user on the standby.

The statistics collector is active during recovery. All scans, reads, blocks, index usage, etc., will be recorded normally on the standby. Replayed actions will not duplicate their effects on primary, so replaying an insert will not increment the Inserts column of pg_stat_user_tables. The stats file is deleted at the start of recovery, so stats from primary and standby will differ; this is considered a feature, not a bug.

Autovacuum is not active during recovery. It will start normally at the end of recovery.

The background writer is active during recovery and will perform restartpoints (similar to checkpoints on the primary) and normal block cleaning activities. This can include updates of the hint bit information stored on the standby server. The CHECKPOINT command is accepted during recovery, though it performs a restartpoint rather than a new checkpoint.
26.5.4. Hot Standby Parameter Reference

Various parameters have been mentioned above in Section 26.5.2 and Section 26.5.3.

On the primary, parameters `wal_level` and `vacuum_defer_cleanup_age` can be used. `max_standby_archive_delay` and `max_standby_streaming_delay` have no effect if set on the primary.

On the standby, parameters `hot_standby`, `max_standby_archive_delay` and `max_standby_streaming_delay` can be used. `vacuum_defer_cleanup_age` has no effect as long as the server remains in standby mode, though it will become relevant if the standby becomes primary.

26.5.5. Caveats

There are several limitations of Hot Standby. These can and probably will be fixed in future releases:

- Operations on hash indexes are not presently WAL-logged, so replay will not update these indexes.
- Full knowledge of running transactions is required before snapshots can be taken. Transactions that use large numbers of subtransactions (currently greater than 64) will delay the start of read only connections until the completion of the longest running write transaction. If this situation occurs, explanatory messages will be sent to the server log.
- Valid starting points for standby queries are generated at each checkpoint on the master. If the standby is shut down while the master is in a shutdown state, it might not be possible to re-enter Hot Standby until the primary is started up, so that it generates further starting points in the WAL logs. This situation isn’t a problem in the most common situations where it might happen. Generally, if the primary is shut down and not available anymore, that’s likely due to a serious failure that requires the standby being converted to operate as the new primary anyway. And in situations where the primary is being intentionally taken down, coordinating to make sure the standby becomes the new primary smoothly is also standard procedure.
- At the end of recovery, `AccessExclusiveLocks` held by prepared transactions will require twice the normal number of lock table entries. If you plan on running either a large number of concurrent prepared transactions that normally take `AccessExclusiveLocks`, or you plan on having one large transaction that takes many `AccessExclusiveLocks`, you are advised to select a larger value of `max_locks_per_transaction`, perhaps as much as twice the value of the parameter on the primary server. You need not consider this at all if your setting of `max_prepared_transactions` is 0.
- The Serializable transaction isolation level is not yet available in hot standby. (See Section 13.2.3 and Section 13.4.1 for details.) An attempt to set a transaction to the serializable isolation level in hot standby mode will generate an error.
Chapter 27. Recovery Configuration

This chapter describes the settings available in the recovery.conf file. They apply only for the duration of the recovery. They must be reset for any subsequent recovery you wish to perform. They cannot be changed once recovery has begun.

Settings in recovery.conf are specified in the format name = 'value'. One parameter is specified per line. Hash marks (#) designate the rest of the line as a comment. To embed a single quote in a parameter value, write two quotes (")

A sample file, share/recovery.conf.sample, is provided in the installation’s share/ directory.

27.1. Archive Recovery Settings

restore_command (string)

The local shell command to execute to retrieve an archived segment of the WAL file series. This parameter is required for archive recovery, but optional for streaming replication. Any %f in the string is replaced by the name of the file to retrieve from the archive, and any %p is replaced by the copy destination path name on the server. (The path name is relative to the current working directory, i.e., the cluster's data directory.) Any %r is replaced by the name of the file containing the last valid restart point. That is the earliest file that must be kept to allow a restore to be restartable, so this information can be used to truncate the archive to just the minimum required to support restarting from the current restore. %r is typically only used by warm-standby configurations (see Section 26.2). Write %% to embed an actual % character.

An exception is that if the command was terminated by a signal (other than SIGTERM, which is used as part of a database server shutdown) or an error by the shell (such as command not found), then recovery will abort and the server will not start up.

archive_cleanup_command (string)

This optional parameter specifies a shell command that will be executed at every restartpoint. The purpose of archive_cleanup_command is to provide a mechanism for cleaning up old archived WAL files that are no longer needed by the standby server. Any %r is replaced by the name of the file containing the last valid restart point. That is the earliest file that must be kept to allow a restore to be restartable, and so all files earlier than %r may be safely removed. This information can be used to truncate the archive to just the minimum required to support restart from the current restore. The pg_archivecleanup module is often used in archive_cleanup_command for single-standby configurations, for example:

archive_cleanup_command = 'pg_archivecleanup /mnt/server/archivedir %r'

Note however that if multiple standby servers are restoring from the same archive directory, you will need to ensure that you do not delete WAL files until they are no longer needed by any of the servers. archive_cleanup_command would typically be used in a warm-standby configuration (see Section 26.2). Write %% to embed an actual % character in the command.
If the command returns a nonzero exit status then a warning log message will be written. An exception is that if the command was terminated by a signal or an error by the shell (such as command not found), a fatal error will be raised.

\texttt{recovery\_end\_command(string)}

This parameter specifies a shell command that will be executed once only at the end of recovery. This parameter is optional. The purpose of the \texttt{recovery\_end\_command} is to provide a mechanism for cleanup following replication or recovery. Any \%r is replaced by the name of the file containing the last valid restart point, like in \texttt{archive\_cleanup\_command}.

If the command returns a nonzero exit status then a warning log message will be written and the database will proceed to start up anyway. An exception is that if the command was terminated by a signal or an error by the shell (such as command not found), the database will not proceed with startup.

### 27.2. Recovery Target Settings

By default, recovery will recover to the end of the WAL log. The following parameters can be used to specify an earlier stopping point. At most one of \texttt{recovery\_target}, \texttt{recovery\_target\_name}, \texttt{recovery\_target\_time}, or \texttt{recovery\_target\_xid} can be used; if more than one of these is specified in the configuration file, the last entry will be used.

\texttt{recovery\_target = ‘immediate’}

This parameter specifies that recovery should end as soon as a consistent state is reached, i.e. as early as possible. When restoring from an online backup, this means the point where taking the backup ended.

Technically, this is a string parameter, but ‘immediate’ is currently the only allowed value.

\texttt{recovery\_target\_name(string)}

This parameter specifies the named restore point (created with \texttt{pg\_create\_restore\_point()}) to which recovery will proceed.

\texttt{recovery\_target\_time(timestamp)}

This parameter specifies the time stamp up to which recovery will proceed. The precise stopping point is also influenced by \texttt{recovery\_target\_inclusive}.

\texttt{recovery\_target\_xid(string)}

This parameter specifies the transaction ID up to which recovery will proceed. Keep in mind that while transaction IDs are assigned sequentially at transaction start, transactions can complete in a different numeric order. The transactions that will be recovered are those that committed before (and optionally including) the specified one. The precise stopping point is also influenced by \texttt{recovery\_target\_inclusive}.

The following options further specify the recovery target, and affect what happens when the target is reached:

\texttt{recovery\_target\_inclusive(boolean)}

Specifies whether to stop just after the specified recovery target (\texttt{true}), or just before the recovery target (\texttt{false}). Applies when either \texttt{recovery\_target\_time} or \texttt{recovery\_target\_xid} is specified. This setting controls whether transactions having exactly the target commit time or ID, respectively, will be included in the recovery. Default is \texttt{true}. 

667
recovery_target_timeline (string)

Specifies recovering into a particular timeline. The default is to recover along the same timeline that was current when the base backup was taken. Setting this to latest recovers to the latest timeline found in the archive, which is useful in a standby server. Other than that you only need to set this parameter in complex re-recovery situations, where you need to return to a state that itself was reached after a point-in-time recovery. See Section 25.3.5 for discussion.

recovery_target_action (enum)

Specifies what action the server should take once the recovery target is reached. The default is pause, which means recovery will be paused. promote means the recovery process will finish and the server will start to accept connections. Finally shutdown will stop the server after reaching the recovery target.

The intended use of the pause setting is to allow queries to be executed against the database to check if this recovery target is the most desirable point for recovery. The paused state can be resumed by using pg_xlog_replay_resume() (see Table 9-80), which then causes recovery to end. If this recovery target is not the desired stopping point, then shut down the server, change the recovery target settings to a later target and restart to continue recovery.

The shutdown setting is useful to have the instance ready at the exact replay point desired. The instance will still be able to replay more WAL records (and in fact will have to replay WAL records since the last checkpoint next time it is started).

Note that because recovery.conf will not be renamed when recovery_target_action is set to shutdown, any subsequent start will end with immediate shutdown unless the configuration is changed or the recovery.conf file is removed manually.

This setting has no effect if no recovery target is set. If hot_standby is not enabled, a setting of pause will act the same as shutdown.

27.3. Standby Server Settings

standby_mode (boolean)

Specifies whether to start the PostgreSQL server as a standby. If this parameter is on, the server will not stop recovery when the end of archived WAL is reached, but will keep trying to continue recovery by fetching new WAL segments using restore_command and/or by connecting to the primary server as specified by the primary_conninfo setting.

primary_conninfo (string)

Specifies a connection string to be used for the standby server to connect with the primary. This string is in the format described in Section 32.1.1. If any option is unspecified in this string, then the corresponding environment variable (see Section 32.14) is checked. If the environment variable is not set either, then defaults are used.

The connection string should specify the host name (or address) of the primary server, as well as the port number if it is not the same as the standby server’s default. Also specify a user name corresponding to a suitably-privileged role on the primary (see Section 26.2.5.1). A password needs to be provided too, if the primary demands password authentication. It can be provided in the primary_conninfo string, or in a separate ~/.pgpass file on the standby server (use replication as the database name). Do not specify a database name in the primary_conninfo string.

This setting has no effect if standby_mode is off.
primary_slot_name (string)

Optionally specifies an existing replication slot to be used when connecting to the primary via streaming replication to control resource removal on the upstream node (see Section 26.2.6). This setting has no effect if primary_conninfo is not set.

trigger_file (string)

Specifies a trigger file whose presence ends recovery in the standby. Even if this value is not set, you can still promote the standby using pg_ctl promote. This setting has no effect if standby_mode is off.

recovery_min_apply_delay (integer)

By default, a standby server restores WAL records from the primary as soon as possible. It may be useful to have a time-delayed copy of the data, offering opportunities to correct data loss errors. This parameter allows you to delay recovery by a fixed period of time, measured in milliseconds if no unit is specified. For example, if you set this parameter to 5min, the standby will replay each transaction commit only when the system time on the standby is at least five minutes past the commit time reported by the master.

It is possible that the replication delay between servers exceeds the value of this parameter, in which case no delay is added. Note that the delay is calculated between the WAL time stamp as written on master and the current time on the standby. Delays in transfer because of network lag or cascading replication configurations may reduce the actual wait time significantly. If the system clocks on master and standby are not synchronized, this may lead to recovery applying records earlier than expected; but that is not a major issue because useful settings of this parameter are much larger than typical time deviations between servers.

The delay occurs only on WAL records for transaction commits. Other records are replayed as quickly as possible, which is not a problem because MVCC visibility rules ensure their effects are not visible until the corresponding commit record is applied.

The delay occurs once the database in recovery has reached a consistent state, until the standby is promoted or triggered. After that the standby will end recovery without further waiting.

This parameter is intended for use with streaming replication deployments; however, if the parameter is specified it will be honored in all cases. hot_standby_feedback will be delayed by use of this feature which could lead to bloat on the master; use both together with care.

Warning

Synchronous replication is affected by this setting when synchronous_commit is set to remote_apply; every COMMIT will need to wait to be applied.
Chapter 28. Monitoring Database Activity

A database administrator frequently wonders, “What is the system doing right now?” This chapter discusses how to find that out.

Several tools are available for monitoring database activity and analyzing performance. Most of this chapter is devoted to describing PostgreSQL’s statistics collector, but one should not neglect regular Unix monitoring programs such as `ps`, `top`, `iostat`, and `vmstat`. Also, once one has identified a poorly-performing query, further investigation might be needed using PostgreSQL’s EXPLAIN command. Section 14.1 discusses EXPLAIN and other methods for understanding the behavior of an individual query.

28.1. Standard Unix Tools

On most Unix platforms, PostgreSQL modifies its command title as reported by `ps`, so that individual server processes can readily be identified. A sample display is

```
$ ps auxww | grep ^postgres
postgres 15551 0.0 0.1 57536 7132 pts/0 S 18:02 0:00 postgres -i
postgres 15554 0.0 0.0 57536 1184 ? Ss 18:02 0:00 postgres: writer process
postgres 15555 0.0 0.0 57536 916 ? Ss 18:02 0:00 postgres: checkpointer process
postgres 15556 0.0 0.0 57536 916 ? Ss 18:02 0:00 postgres: wal writer process
postgres 15557 0.0 0.0 58504 2244 ? Ss 18:02 0:00 postgres: autovacuum launcher process
postgres 15558 0.0 0.0 17512 1068 ? Ss 18:02 0:00 postgres: stats collector process
postgres 15582 0.0 0.0 58772 3080 ? Ss 18:04 0:00 postgres: joe runbug 127.0.0.1 idle
postgres 15606 0.0 0.0 58772 3052 ? Ss 18:07 0:00 postgres: tgl regression [local] idle in transaction
poststats 15610 0.0 0.0 58772 3056 ? Ss 18:07 0:00 postgres: tgl regression
```

(The appropriate invocation of `ps` varies across different platforms, as do the details of what is shown. This example is from a recent Linux system.) The first process listed here is the master server process. The command arguments shown for it are the same ones used when it was launched. The next five processes are background worker processes automatically launched by the master process. (The “stats collector” process will not be present if you have set the system not to start the statistics collector; likewise the “autovacuum launcher” process can be disabled.) Each of the remaining processes is a server process handling one client connection. Each such process sets its command line display in the form

```
postgres: user database host activity
```

The user, database, and (client) host items remain the same for the life of the client connection, but the activity indicator changes. The activity can be `idle` (i.e., waiting for a client command), `idle in transaction` (waiting for client inside a `BEGIN` block), or a command type name such as `SELECT`. Also, `waiting` is appended if the server process is presently waiting on a lock held by another session. In the above example we can infer that process 15606 is waiting for process 15610 to complete its transaction and thereby release some lock. (Process 15610 must be the blocker, because there is no other active session. In more complicated cases it would be necessary to look into the `pg_locks` system view to determine who is blocking whom.)

If `cluster_name` has been configured the cluster name will also be shown in `ps` output:

```
$ psql -c 'SHOW cluster_name'
cluster_name
-----------------
server1
```

670
Chapter 28. Monitoring Database Activity

28.2. The Statistics Collector

PostgreSQL’s statistics collector is a subsystem that supports collection and reporting of information about server activity. Presently, the collector can count accesses to tables and indexes in both disk-block and individual-row terms. It also tracks the total number of rows in each table, and information about vacuum and analyze actions for each table. It can also count calls to user-defined functions and the total time spent in each one.

PostgreSQL also supports reporting dynamic information about exactly what is going on in the system right now, such as the exact command currently being executed by other server processes, and which other connections exist in the system. This facility is independent of the collector process.

28.2.1. Statistics Collection Configuration

Since collection of statistics adds some overhead to query execution, the system can be configured to collect or not collect information. This is controlled by configuration parameters that are normally set in postgresql.conf. (See Chapter 19 for details about setting configuration parameters.)

The parameter track_activities enables monitoring of the current command being executed by any server process.

The parameter track_counts controls whether statistics are collected about table and index accesses.

The parameter track_functions enables tracking of usage of user-defined functions.

The parameter track_io_timing enables monitoring of block read and write times.

Normally these parameters are set in postgresql.conf so that they apply to all server processes, but it is possible to turn them on or off in individual sessions using the SET command. (To prevent ordinary users from hiding their activity from the administrator, only superusers are allowed to change these parameters with SET.)

The statistics collector transmits the collected information to other PostgreSQL processes through temporary files. These files are stored in the directory named by the stats_temp_directory parameter, pg_stat_tmp by default. For better performance, stats_temp_directory can be pointed at a RAM-based file system, decreasing physical I/O requirements. When the server shuts down cleanly,
a permanent copy of the statistics data is stored in the pg_stat subdirectory, so that statistics can be retained across server restarts. When recovery is performed at server start (e.g. after immediate shutdown, server crash, and point-in-time recovery), all statistics counters are reset.

### 28.2.2. Viewing Statistics

Several predefined views, listed in Table 28-1, are available to show the current state of the system. There are also several other views, listed in Table 28-2, available to show the results of statistics collection. Alternatively, one can build custom views using the underlying statistics functions, as discussed in Section 28.2.3.

When using the statistics to monitor collected data, it is important to realize that the information does not update instantaneously. Each individual server process transmits new statistical counts to the collector just before going idle; so a query or transaction still in progress does not affect the displayed totals. Also, the collector itself emits a new report at most once per `PGSTAT_STAT_INTERVAL` milliseconds (500 ms unless altered while building the server). So the displayed information lags behind actual activity. However, current-query information collected by `track_activities` is always up-to-date.

Another important point is that when a server process is asked to display any of these statistics, it first fetches the most recent report emitted by the collector process and then continues to use this snapshot for all statistical views and functions until the end of its current transaction. So the statistics will show static information as long as you continue the current transaction. Similarly, information about the current queries of all sessions is collected when any such information is first requested within a transaction, and the same information will be displayed throughout the transaction. This is a feature, not a bug, because it allows you to perform several queries on the statistics and correlate the results without worrying that the numbers are changing underneath you. But if you want to see new results with each query, be sure to do the queries outside any transaction block. Alternatively, you can invoke `pg_stat_clear_snapshot()`, which will discard the current transaction’s statistics snapshot (if any). The next use of statistical information will cause a new snapshot to be fetched.

A transaction can also see its own statistics (as yet untransmitted to the collector) in the views `pg_stat_xact_all_tables`, `pg_stat_xact_sys_tables`, `pg_stat_xact_user_tables`, and `pg_stat_xact_user_functions`. These numbers do not act as stated above; instead they update continuously throughout the transaction.

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_stat_activity</code></td>
<td>One row per server process, showing information related to the current activity of that process, such as state and current query. See <code>pg_stat_activity</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_replication</code></td>
<td>One row per WAL sender process, showing statistics about replication to that sender’s connected standby server. See <code>pg_stat_replication</code> for details.</td>
</tr>
<tr>
<td><code>pg_stat_wal_receiver</code></td>
<td>Only one row, showing statistics about the WAL receiver from that receiver’s connected server. See <code>pg_stat_wal_receiver</code> for details.</td>
</tr>
</tbody>
</table>
### Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_stat_ssl</td>
<td>One row per connection (regular and replication), showing information about SSL used on this connection. See pg_stat_ssl for details.</td>
</tr>
<tr>
<td>pg_stat_progress_vacuum</td>
<td>One row for each backend (including autovacuum worker processes) running VACUUM, showing current progress. See Section 28.4.1.</td>
</tr>
</tbody>
</table>

**Table 28-2. Collected Statistics Views**

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_stat_archiver</td>
<td>One row only, showing statistics about the WAL archiver process’s activity. See pg_stat_archiver for details.</td>
</tr>
<tr>
<td>pg_stat_bgwriter</td>
<td>One row only, showing statistics about the background writer process’s activity. See pg_stat_bgwriter for details.</td>
</tr>
<tr>
<td>pg_stat_database</td>
<td>One row per database, showing database-wide statistics. See pg_stat_database for details.</td>
</tr>
<tr>
<td>pg_stat_database_conflicts</td>
<td>One row per database, showing database-wide statistics about query cancels due to conflict with recovery on standby servers. See pg_stat_database_conflicts for details.</td>
</tr>
<tr>
<td>pg_stat_all_tables</td>
<td>One row for each table in the current database, showing statistics about accesses to that specific table. See pg_stat_all_tables for details.</td>
</tr>
<tr>
<td>pg_stat_sys_tables</td>
<td>Same as pg_stat_all_tables, except that only system tables are shown.</td>
</tr>
<tr>
<td>pg_stat_user_tables</td>
<td>Same as pg_stat_all_tables, except that only user tables are shown.</td>
</tr>
<tr>
<td>pg_stat_xact_all_tables</td>
<td>Similar to pg_stat_all_tables, but counts actions taken so far within the current transaction (which are not yet included in pg_stat_all_tables and related views). The columns for numbers of live and dead rows and vacuum and analyze actions are not present in this view.</td>
</tr>
<tr>
<td>pg_stat_xact_sys_tables</td>
<td>Same as pg_stat_xact_all_tables, except that only system tables are shown.</td>
</tr>
<tr>
<td>pg_stat_xact_user_tables</td>
<td>Same as pg_stat_xact_all_tables, except that only user tables are shown.</td>
</tr>
<tr>
<td>pg_stat_all_indexes</td>
<td>One row for each index in the current database, showing statistics about accesses to that specific index. See pg_stat_all_indexes for details.</td>
</tr>
<tr>
<td>pg_stat_sys_indexes</td>
<td>Same as pg_stat_all_indexes, except that only indexes on system tables are shown.</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>View Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_stat_user_indexes</td>
<td>Same as pg_stat_all_indexes, except that only indexes on user tables are shown.</td>
</tr>
<tr>
<td>pg_statio_all_tables</td>
<td>One row for each table in the current database, showing statistics about I/O on that specific table. See pg_statio_all_tables for details.</td>
</tr>
<tr>
<td>pg_statio_sys_tables</td>
<td>Same as pg_statio_all_tables, except that only system tables are shown.</td>
</tr>
<tr>
<td>pg_statio_user_tables</td>
<td>Same as pg_statio_all_tables, except that only user tables are shown.</td>
</tr>
<tr>
<td>pg_statio_all_indexes</td>
<td>One row for each index in the current database, showing statistics about I/O on that specific index. See pg_statio_all_indexes for details.</td>
</tr>
<tr>
<td>pg_statio_sys_indexes</td>
<td>Same as pg_statio_all_indexes, except that only indexes on system tables are shown.</td>
</tr>
<tr>
<td>pg_statio_user_indexes</td>
<td>Same as pg_statio_all_indexes, except that only indexes on user tables are shown.</td>
</tr>
<tr>
<td>pg_statio_all_sequences</td>
<td>One row for each sequence in the current database, showing statistics about I/O on that specific sequence. See pg_statio_all_sequences for details.</td>
</tr>
<tr>
<td>pg_statio_sys_sequences</td>
<td>Same as pg_statio_all_sequences, except that only system sequences are shown. (Presently, no system sequences are defined, so this view is always empty.)</td>
</tr>
<tr>
<td>pg_statio_user_sequences</td>
<td>Same as pg_statio_all_sequences, except that only user sequences are shown.</td>
</tr>
<tr>
<td>pg_stat_user_functions</td>
<td>One row for each tracked function, showing statistics about executions of that function. See pg_stat_user_functions for details.</td>
</tr>
<tr>
<td>pg_stat_xact_user_functions</td>
<td>Similar to pg_stat_user_functions, but counts only calls during the current transaction (which are not yet included in pg_stat_user_functions).</td>
</tr>
</tbody>
</table>

The per-index statistics are particularly useful to determine which indexes are being used and how effective they are.

The pg_statio_ views are primarily useful to determine the effectiveness of the buffer cache. When the number of actual disk reads is much smaller than the number of buffer hits, then the cache is satisfying most read requests without invoking a kernel call. However, these statistics do not give the entire story: due to the way in which PostgreSQL handles disk I/O, data that is not in the PostgreSQL buffer cache might still reside in the kernel’s I/O cache, and might therefore still be fetched without requiring a physical read. Users interested in obtaining more detailed information on PostgreSQL I/O behavior are advised to use the PostgreSQL statistics collector in combination with operating system utilities that allow insight into the kernel’s handling of I/O.

Table 28-3. pg_stat_activity View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
</table>

674
### Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of the database this backend is connected to</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of the database this backend is connected to</td>
</tr>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of this backend</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td>OID of the user logged into this backend</td>
</tr>
<tr>
<td>usename</td>
<td>name</td>
<td>Name of the user logged into this backend</td>
</tr>
<tr>
<td>application_name</td>
<td>text</td>
<td>Name of the application that is connected to this backend</td>
</tr>
<tr>
<td>client_addr</td>
<td>inet</td>
<td>IP address of the client connected to this backend. If this field is null, it indicates either that the client is connected via a Unix socket on the server machine or that this is an internal process such as autovacuum.</td>
</tr>
<tr>
<td>client_hostname</td>
<td>text</td>
<td>Host name of the connected client, as reported by a reverse DNS lookup of client_addr. This field will only be non-null for IP connections, and only when log_hostname is enabled.</td>
</tr>
<tr>
<td>client_port</td>
<td>integer</td>
<td>TCP port number that the client is using for communication with this backend, or -1 if a Unix socket is used</td>
</tr>
<tr>
<td>backend_start</td>
<td>timestamp with time zone</td>
<td>Time when this process was started, i.e., when the client connected to the server</td>
</tr>
<tr>
<td>xact_start</td>
<td>timestamp with time zone</td>
<td>Time when this process’ current transaction was started, or null if no transaction is active. If the current query is the first of its transaction, this column is equal to the query_start column.</td>
</tr>
<tr>
<td>query_start</td>
<td>timestamp with time zone</td>
<td>Time when the currently active query was started, or if state is not active, when the last query was started</td>
</tr>
<tr>
<td>state_change</td>
<td>timestamp with time zone</td>
<td>Time when the state was last changed</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>--------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>wait_event_type</td>
<td>text</td>
<td>The type of event for which the backend is waiting, if any; otherwise NULL. Possible values are: • LWLockNamed: The backend is waiting for a specific named lightweight lock. Each such lock protects a particular data structure in shared memory. <em>wait_event</em> will contain the name of the lightweight lock.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• LWLockTranche: The backend is waiting for one of a group of related lightweight locks. All locks in the group perform a similar function; <em>wait_event</em> will identify the general purpose of locks in that group.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• Lock: The backend is waiting for a heavyweight lock. Heavyweight locks, also known as lock manager locks or simply locks, primarily protect SQL-visible objects such as tables. However, they are also used to ensure mutual exclusion for certain internal operations such as relation extension. <em>wait_event</em> will identify the type of lock awaited.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• BufferPin: The server process is waiting to access a data buffer during a period when no other process can be examining that buffer. Buffer pin waits can be protracted if another process holds an open cursor which last read data from the buffer in question.</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>wait_event</td>
<td>text</td>
<td>Wait event name if backend is currently waiting, otherwise NULL. See Table 28-4 for details.</td>
</tr>
<tr>
<td>state</td>
<td>text</td>
<td>Current overall state of this backend. Possible values are:</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• active: The backend is executing a query.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• idle: The backend is waiting for a new client command.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• idle in transaction: The backend is in a transaction, but is not currently executing a query.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• idle in transaction (aborted): This state is similar to idle in transaction, except one of the statements in the transaction caused an error.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• fastpath function call: The backend is executing a fast-path function.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>• disabled: This state is reported if track_activities is disabled in this backend.</td>
</tr>
<tr>
<td>backend_xid</td>
<td>xid</td>
<td>Top-level transaction identifier of this backend, if any.</td>
</tr>
<tr>
<td>backend_xmin</td>
<td>xid</td>
<td>The current backend’s xmin horizon.</td>
</tr>
<tr>
<td>query</td>
<td>text</td>
<td>Text of this backend’s most recent query. If state is active this field shows the currently executing query. In all other states, it shows the last query that was executed.</td>
</tr>
</tbody>
</table>

The pg_stat_activity view will have one row per server process, showing information related to the current activity of that process.

Note: The wait_event and state columns are independent. If a backend is in the active state, it may or may not be waiting on some event. If the state is active and wait_event is non-null,
it means that a query is being executed, but is being blocked somewhere in the system.

<table>
<thead>
<tr>
<th>Wait Event Type</th>
<th>Wait Event Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>LWLockNamed</td>
<td>ShmemIndexLock</td>
<td>Waiting to find or allocate space in shared memory.</td>
</tr>
<tr>
<td></td>
<td>OidGenLock</td>
<td>Waiting to allocate or assign an OID.</td>
</tr>
<tr>
<td></td>
<td>XidGenLock</td>
<td>Waiting to allocate or assign a transaction id.</td>
</tr>
<tr>
<td></td>
<td>ProcArrayLock</td>
<td>Waiting to get a snapshot or clearing a transaction id at transaction end.</td>
</tr>
<tr>
<td></td>
<td>SInvalReadLock</td>
<td>Waiting to retrieve or remove messages from shared invalidation queue.</td>
</tr>
<tr>
<td></td>
<td>SInvalWriteLock</td>
<td>Waiting to add a message in shared invalidation queue.</td>
</tr>
<tr>
<td></td>
<td>WALBufMappingLock</td>
<td>Waiting to replace a page in WAL buffers.</td>
</tr>
<tr>
<td></td>
<td>WALWriteLock</td>
<td>Waiting for WAL buffers to be written to disk.</td>
</tr>
<tr>
<td></td>
<td>ControlFileLock</td>
<td>Waiting to read or update the control file or creation of a new WAL file.</td>
</tr>
<tr>
<td></td>
<td>CheckpointLock</td>
<td>Waiting to perform checkpoint.</td>
</tr>
<tr>
<td></td>
<td>CLogControlLock</td>
<td>Waiting to read or update transaction status.</td>
</tr>
<tr>
<td></td>
<td>SubtransControlLock</td>
<td>Waiting to read or update sub-transaction information.</td>
</tr>
<tr>
<td></td>
<td>MultiXactGenLock</td>
<td>Waiting to read or update shared multixact state.</td>
</tr>
<tr>
<td></td>
<td>MultiXactOffsetControlLock</td>
<td>Waiting to read or update multixact offset mappings.</td>
</tr>
<tr>
<td></td>
<td>MultiXactMemberControlLock</td>
<td>Waiting to read or update multixact member mappings.</td>
</tr>
<tr>
<td></td>
<td>RelCacheInitLock</td>
<td>Waiting to read or write relation cache initialization file.</td>
</tr>
<tr>
<td></td>
<td>CheckpointerCommLock</td>
<td>Waiting to manage fsync requests.</td>
</tr>
<tr>
<td></td>
<td>TwoPhaseStateLock</td>
<td>Waiting to read or update the state of prepared transactions.</td>
</tr>
<tr>
<td></td>
<td>TablespaceCreateLock</td>
<td>Waiting to create or drop the tablespace.</td>
</tr>
<tr>
<td>Wait Event Type</td>
<td>Wait Event Name</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------------</td>
<td>--------------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>BtreeVacuumLock</td>
<td>BtreeVacuumLock</td>
<td>Waiting to read or update vacuum-related information for a B-tree index.</td>
</tr>
<tr>
<td>AddinShmemInitLock</td>
<td>AddinShmemInitLock</td>
<td>Waiting to manage space allocation in shared memory.</td>
</tr>
<tr>
<td>AutovacuumLock</td>
<td>AutovacuumLock</td>
<td>Autovacuum worker or launcher waiting to update or read the current state of autovacuum workers.</td>
</tr>
<tr>
<td>AutovacuumScheduleLock</td>
<td>AutovacuumScheduleLock</td>
<td>Waiting to ensure that the table it has selected for a vacuum still needs vacuuming.</td>
</tr>
<tr>
<td>SyncScanLock</td>
<td>SyncScanLock</td>
<td>Waiting to get the start location of a scan on a table for synchronized scans.</td>
</tr>
<tr>
<td>RelationMappingLock</td>
<td>RelationMappingLock</td>
<td>Waiting to update the relation map file used to store catalog to filenode mapping.</td>
</tr>
<tr>
<td>AsyncCtlLock</td>
<td>AsyncCtlLock</td>
<td>Waiting to read or update shared notification state.</td>
</tr>
<tr>
<td>AsyncQueueLock</td>
<td>AsyncQueueLock</td>
<td>Waiting to read or update notification messages.</td>
</tr>
<tr>
<td>SerializableXactHashLock</td>
<td>SerializableXactHashLock</td>
<td>Waiting to retrieve or store information about serializable transactions.</td>
</tr>
<tr>
<td>SerializableFinishedList</td>
<td>SerializableFinishedList</td>
<td>Waiting to access the list of finished serializable transactions.</td>
</tr>
<tr>
<td>SerializablePredicateLock</td>
<td>SerializablePredicateLock</td>
<td>Waiting to perform an operation on a list of locks held by serializable transactions.</td>
</tr>
<tr>
<td>OldSerXidLock</td>
<td>OldSerXidLock</td>
<td>Waiting to read or record conflicting serializable transactions.</td>
</tr>
<tr>
<td>SyncRepLock</td>
<td>SyncRepLock</td>
<td>Waiting to read or update information about synchronous replicas.</td>
</tr>
<tr>
<td>BackgroundWorkerLock</td>
<td>BackgroundWorkerLock</td>
<td>Waiting to read or update background worker state.</td>
</tr>
<tr>
<td>DynamicSharedMemoryControlLock</td>
<td>DynamicSharedMemoryControlLock</td>
<td>Waiting to read or update dynamic shared memory state.</td>
</tr>
<tr>
<td>AutoFileLock</td>
<td>AutoFileLock</td>
<td>Waiting to update the <code>postgresql.auto.conf</code> file.</td>
</tr>
<tr>
<td>ReplicationSlotAllocation</td>
<td>ReplicationSlotAllocation</td>
<td>Waiting to allocate or free a replication slot.</td>
</tr>
<tr>
<td>ReplicationSlotControlLock</td>
<td>ReplicationSlotControlLock</td>
<td>Waiting to read or update replication slot state.</td>
</tr>
</tbody>
</table>
## Wait Event Type

<table>
<thead>
<tr>
<th>Wait Event Type</th>
<th>Wait Event Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>CommitTsControlLock</td>
<td>Waiting to read or update transaction commit timestamps.</td>
<td></td>
</tr>
<tr>
<td>CommitTsLock</td>
<td>Waiting to read or update the last value set for the transaction timestamp.</td>
<td></td>
</tr>
<tr>
<td>ReplicationOriginLock</td>
<td>Waiting to setup, drop or use replication origin.</td>
<td></td>
</tr>
<tr>
<td>MultiXactTruncationLock</td>
<td>Waiting to read or truncate multixact information.</td>
<td></td>
</tr>
<tr>
<td>OldSnapshotTimeMapLock</td>
<td>Waiting to read or update old snapshot control information.</td>
<td></td>
</tr>
<tr>
<td>LWLockTranche</td>
<td>clog</td>
<td>Waiting for I/O on a clog (transaction status) buffer.</td>
</tr>
<tr>
<td></td>
<td>commit_timestamp</td>
<td>Waiting for I/O on commit timestamp buffer.</td>
</tr>
<tr>
<td></td>
<td>subtrans</td>
<td>Waiting for I/O a subtransaction buffer.</td>
</tr>
<tr>
<td></td>
<td>multixact_offset</td>
<td>Waiting for I/O on a multixact offset buffer.</td>
</tr>
<tr>
<td></td>
<td>multixact_member</td>
<td>Waiting for I/O on a multixact_member buffer.</td>
</tr>
<tr>
<td></td>
<td>async</td>
<td>Waiting for I/O on an async (notify) buffer.</td>
</tr>
<tr>
<td></td>
<td>oldserxid</td>
<td>Waiting to I/O on an oldserxid buffer.</td>
</tr>
<tr>
<td></td>
<td>wal_insert</td>
<td>Waiting to insert WAL into a memory buffer.</td>
</tr>
<tr>
<td></td>
<td>buffer_content</td>
<td>Waiting to read or write a data page in memory.</td>
</tr>
<tr>
<td></td>
<td>buffer_io</td>
<td>Waiting for I/O on a data page.</td>
</tr>
<tr>
<td></td>
<td>replication_origin</td>
<td>Waiting to read or update the replication progress.</td>
</tr>
<tr>
<td></td>
<td>replication_slot_io</td>
<td>Waiting for I/O on a replication slot.</td>
</tr>
<tr>
<td></td>
<td>proc</td>
<td>Waiting to read or update the fast-path lock information.</td>
</tr>
<tr>
<td></td>
<td>buffer_mapping</td>
<td>Waiting to associate a data block with a buffer in the buffer pool.</td>
</tr>
<tr>
<td></td>
<td>lock_manager</td>
<td>Waiting to add or examine locks for backends, or waiting to join or exit a locking group (used by parallel query).</td>
</tr>
<tr>
<td></td>
<td>predicate_lock_manager</td>
<td>Waiting to add or examine predicate lock information.</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Wait Event Type</th>
<th>Wait Event Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Lock</td>
<td>relation</td>
<td>Waiting to acquire a lock on a relation.</td>
</tr>
<tr>
<td>extend</td>
<td></td>
<td>Waiting to extend a relation.</td>
</tr>
<tr>
<td>page</td>
<td></td>
<td>Waiting to acquire a lock on page of a relation.</td>
</tr>
<tr>
<td>tuple</td>
<td></td>
<td>Waiting to acquire a lock on a tuple.</td>
</tr>
<tr>
<td>transactionid</td>
<td></td>
<td>Waiting for a transaction to finish.</td>
</tr>
<tr>
<td>virtualxid</td>
<td></td>
<td>Waiting to acquire a virtual xid lock.</td>
</tr>
<tr>
<td>speculative token</td>
<td></td>
<td>Waiting to acquire a speculative insertion lock.</td>
</tr>
<tr>
<td>object</td>
<td></td>
<td>Waiting to acquire a lock on a non-relation database object.</td>
</tr>
<tr>
<td>userlock</td>
<td></td>
<td>Waiting to acquire a userlock.</td>
</tr>
<tr>
<td>advisory</td>
<td></td>
<td>Waiting to acquire an advisory user lock.</td>
</tr>
</tbody>
</table>

| BufferPin          | BufferPin         | Waiting to acquire a pin on a buffer.            |

Note: For tranches registered by extensions, the name is specified by extension and this will be displayed as wait_event. It is quite possible that user has registered the tranche in one of the backends (by having allocation in dynamic shared memory) in which case other backends won’t have that information, so we display extension for such cases.

Here is an example of how wait events can be viewed

```
SELECT pid, wait_event_type, wait_event FROM pg_stat_activity WHERE wait_event is NOT NULL;
```

<table>
<thead>
<tr>
<th>pid</th>
<th>wait_event_type</th>
<th>wait_event</th>
</tr>
</thead>
<tbody>
<tr>
<td>2540</td>
<td>Lock</td>
<td>relation</td>
</tr>
<tr>
<td>6644</td>
<td>LWLockNamed</td>
<td>ProcArrayLock</td>
</tr>
</tbody>
</table>

(2 rows)

Table 28-5. pg_stat_replication View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of a WAL sender process</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td>OID of the user logged into this WAL sender process</td>
</tr>
<tr>
<td>username</td>
<td>name</td>
<td>Name of the user logged into this WAL sender process</td>
</tr>
</tbody>
</table>
### Column Names and Descriptions

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>application_name</td>
<td>text</td>
<td>Name of the application that is connected to this WAL sender.</td>
</tr>
<tr>
<td>client_addr</td>
<td>inet</td>
<td>IP address of the client connected to this WAL sender. If this field is null, it indicates that the client is connected via a Unix socket on the server machine.</td>
</tr>
<tr>
<td>clientHostname</td>
<td>text</td>
<td>Host name of the connected client, as reported by a reverse DNS lookup of client_addr. This field will only be non-null for IP connections, and only when log_hostname is enabled.</td>
</tr>
<tr>
<td>client_port</td>
<td>integer</td>
<td>TCP port number that the client is using for communication with this WAL sender, or -1 if a Unix socket is used.</td>
</tr>
<tr>
<td>backend_start</td>
<td>timestamp with time zone</td>
<td>Time when this process was started, i.e., when the client connected to this WAL sender</td>
</tr>
<tr>
<td>backend_xmin</td>
<td>xid</td>
<td>This standby’s xmin horizon reported by hot_standby_feedback.</td>
</tr>
<tr>
<td>state</td>
<td>text</td>
<td>Current WAL sender state</td>
</tr>
<tr>
<td>sent_location</td>
<td>pg_lsn</td>
<td>Last transaction log position sent on this connection</td>
</tr>
<tr>
<td>write_location</td>
<td>pg_lsn</td>
<td>Last transaction log position written to disk by this standby server</td>
</tr>
<tr>
<td>flush_location</td>
<td>pg_lsn</td>
<td>Last transaction log position flushed to disk by this standby server</td>
</tr>
<tr>
<td>replay_location</td>
<td>pg_lsn</td>
<td>Last transaction log position replayed into the database on this standby server</td>
</tr>
<tr>
<td>sync_priority</td>
<td>integer</td>
<td>Priority of this standby server for being chosen as the synchronous standby</td>
</tr>
<tr>
<td>sync_state</td>
<td>text</td>
<td>Synchronous state of this standby server</td>
</tr>
</tbody>
</table>

The `pg_stat_replication` view will contain one row per WAL sender process, showing statistics about replication to that sender’s connected standby server. Only directly connected standbys are listed; no information is available about downstream standby servers.
Table 28-6. pg_stat_wal_receiver View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of the WAL receiver process</td>
</tr>
<tr>
<td>status</td>
<td>text</td>
<td>Activity status of the WAL receiver process</td>
</tr>
<tr>
<td>receive_start_lsn</td>
<td>pg_lsn</td>
<td>First transaction log position used when WAL receiver is started</td>
</tr>
<tr>
<td>receive_start_tli</td>
<td>integer</td>
<td>First timeline number used when WAL receiver is started</td>
</tr>
<tr>
<td>received_lsn</td>
<td>pg_lsn</td>
<td>Last transaction log position already received and flushed to disk, the initial value of this field being the first log position used when WAL receiver is started</td>
</tr>
<tr>
<td>received_tli</td>
<td>integer</td>
<td>Timeline number of last transaction log position received and flushed to disk, the initial value of this field being the timeline number of the first log position used when WAL receiver is started</td>
</tr>
<tr>
<td>last_msg_send_time</td>
<td>timestamp with time zone</td>
<td>Send time of last message received from origin WAL sender</td>
</tr>
<tr>
<td>last_msg_receipt_time</td>
<td>timestamp with time zone</td>
<td>Receipt time of last message received from origin WAL sender</td>
</tr>
<tr>
<td>latest_end_lsn</td>
<td>pg_lsn</td>
<td>Last transaction log position reported to origin WAL sender</td>
</tr>
<tr>
<td>latest_end_time</td>
<td>timestamp with time zone</td>
<td>Time of last transaction log position reported to origin WAL sender</td>
</tr>
<tr>
<td>slot_name</td>
<td>text</td>
<td>Replication slot name used by this WAL receiver</td>
</tr>
<tr>
<td>conninfo</td>
<td>text</td>
<td>Connection string used by this WAL receiver, with security-sensitive fields obfuscated.</td>
</tr>
</tbody>
</table>

The pg_stat_wal_receiver view will contain only one row, showing statistics about the WAL receiver from that receiver’s connected server.

Table 28-7. pg_stat_ssl View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
</table>

683
Chapter 28. Monitoring Database Activity

### Table 28-8. pg_stat_ssl View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of a backend or WAL sender process</td>
</tr>
<tr>
<td>ssl</td>
<td>boolean</td>
<td>True if SSL is used on this connection</td>
</tr>
<tr>
<td>version</td>
<td>text</td>
<td>Version of SSL in use, or NULL if SSL is not in use on this connection</td>
</tr>
<tr>
<td>cipher</td>
<td>text</td>
<td>Name of SSL cipher in use, or NULL if SSL is not in use on this connection</td>
</tr>
<tr>
<td>bits</td>
<td>integer</td>
<td>Number of bits in the encryption algorithm used, or NULL if SSL is not used on this connection</td>
</tr>
<tr>
<td>compression</td>
<td>boolean</td>
<td>True if SSL compression is in use, false if not, or NULL if SSL is not in use on this connection</td>
</tr>
<tr>
<td>clientdn</td>
<td>text</td>
<td>Distinguished Name (DN) field from the client certificate used, or NULL if no client certificate was supplied or if SSL is not in use on this connection. This field is truncated if the DN field is longer than NAMEDATALEN (64 characters in a standard build)</td>
</tr>
</tbody>
</table>

The `pg_stat_ssl` view will contain one row per backend or WAL sender process, showing statistics about SSL usage on this connection. It can be joined to `pg_stat_activity` or `pg_stat_replication` on the `pid` column to get more details about the connection.

### Table 28-8. pg_stat_archiver View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>archived_count</td>
<td>bigint</td>
<td>Number of WAL files that have been successfully archived</td>
</tr>
<tr>
<td>last_archived_wal</td>
<td>text</td>
<td>Name of the last WAL file successfully archived</td>
</tr>
<tr>
<td>last_archived_time</td>
<td>timestamp with time zone</td>
<td>Time of the last successful archive operation</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>stats_reset</td>
<td>timestamp with time zone</td>
<td>Time at which these statistics were last reset</td>
</tr>
</tbody>
</table>

The `pg_stat_archiver` view will always have a single row, containing data about the archiver process of the cluster.

Table 28-9. `pg_stat_bgwriter` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>checkpoints_timed</td>
<td>bigint</td>
<td>Number of scheduled checkpoints that have been performed</td>
</tr>
<tr>
<td>checkpoints_req</td>
<td>bigint</td>
<td>Number of requested checkpoints that have been performed</td>
</tr>
<tr>
<td>checkpoint_write_time</td>
<td>double precision</td>
<td>Total amount of time that has been spent in the portion of checkpoint processing where files are written to disk, in milliseconds</td>
</tr>
<tr>
<td>checkpoint_sync_time</td>
<td>double precision</td>
<td>Total amount of time that has been spent in the portion of checkpoint processing where files are synchronized to disk, in milliseconds</td>
</tr>
<tr>
<td>buffers_checkpoint</td>
<td>bigint</td>
<td>Number of buffers written during checkpoints</td>
</tr>
<tr>
<td>buffers_clean</td>
<td>bigint</td>
<td>Number of buffers written by the background writer</td>
</tr>
<tr>
<td>maxwritten_clean</td>
<td>bigint</td>
<td>Number of times the background writer stopped a cleaning scan because it had written too many buffers</td>
</tr>
<tr>
<td>buffers_backend</td>
<td>bigint</td>
<td>Number of buffers written directly by a backend</td>
</tr>
<tr>
<td>buffers_backend_fsync</td>
<td>bigint</td>
<td>Number of times a backend had to execute its own fsync call (normally the background writer handles those even when the backend does its own write)</td>
</tr>
<tr>
<td>buffers_alloc</td>
<td>bigint</td>
<td>Number of buffers allocated</td>
</tr>
<tr>
<td>stats_reset</td>
<td>timestamp with time zone</td>
<td>Time at which these statistics were last reset</td>
</tr>
</tbody>
</table>

The `pg_stat_bgwriter` view will always have a single row, containing global data for the cluster.

Table 28-10. `pg_stat_database` View
<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of a database</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of this database</td>
</tr>
<tr>
<td>numbackends</td>
<td>integer</td>
<td>Number of backends currently connected to this database. This is the only column in this view that returns a value reflecting current state; all other columns return the accumulated values since the last reset.</td>
</tr>
<tr>
<td>xact_commit</td>
<td>bigint</td>
<td>Number of transactions in this database that have been committed</td>
</tr>
<tr>
<td>xact_rollback</td>
<td>bigint</td>
<td>Number of transactions in this database that have been rolled back</td>
</tr>
<tr>
<td>blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read in this database</td>
</tr>
<tr>
<td>blks_hit</td>
<td>bigint</td>
<td>Number of times disk blocks were found already in the buffer cache, so that a read was not necessary (this only includes hits in the PostgreSQL buffer cache, not the operating system’s file system cache)</td>
</tr>
<tr>
<td>tup_returned</td>
<td>bigint</td>
<td>Number of rows returned by queries in this database</td>
</tr>
<tr>
<td>tup_fetched</td>
<td>bigint</td>
<td>Number of rows fetched by queries in this database</td>
</tr>
<tr>
<td>tup_inserted</td>
<td>bigint</td>
<td>Number of rows inserted by queries in this database</td>
</tr>
<tr>
<td>tup_updated</td>
<td>bigint</td>
<td>Number of rows updated by queries in this database</td>
</tr>
<tr>
<td>tup_deleted</td>
<td>bigint</td>
<td>Number of rows deleted by queries in this database</td>
</tr>
<tr>
<td>conflicts</td>
<td>bigint</td>
<td>Number of queries canceled due to conflicts with recovery in this database. (Conflicts occur only on standby servers; see pg_stat_database_conflicts for details.)</td>
</tr>
</tbody>
</table>
## Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>temp_files</td>
<td>bigint</td>
<td>Number of temporary files created by queries in this database. All temporary files are counted, regardless of why the temporary file was created (e.g., sorting or hashing), and regardless of the log_temp_files setting.</td>
</tr>
<tr>
<td>temp_bytes</td>
<td>bigint</td>
<td>Total amount of data written to temporary files by queries in this database. All temporary files are counted, regardless of why the temporary file was created, and regardless of the log_temp_files setting.</td>
</tr>
<tr>
<td>deadlocks</td>
<td>bigint</td>
<td>Number of deadlocks detected in this database.</td>
</tr>
<tr>
<td>blk_read_time</td>
<td>double precision</td>
<td>Time spent reading data file blocks by backends in this database, in milliseconds.</td>
</tr>
<tr>
<td>blk_write_time</td>
<td>double precision</td>
<td>Time spent writing data file blocks by backends in this database, in milliseconds.</td>
</tr>
<tr>
<td>stats_reset</td>
<td>timestamp with time zone</td>
<td>Time at which these statistics were last reset.</td>
</tr>
</tbody>
</table>

The `pg_stat_database` view will contain one row for each database in the cluster, showing database-wide statistics.

### Table 28-11. `pg_stat_database_conflicts` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of a database</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of this database</td>
</tr>
<tr>
<td>confl_tablespace</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to dropped table spaces</td>
</tr>
<tr>
<td>confl_lock</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to lock timeouts</td>
</tr>
<tr>
<td>confl_snapshot</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to old snapshots</td>
</tr>
<tr>
<td>confl_bufferpin</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to pinned buffers</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

The `pg_stat_database_conflicts` view will contain one row per database, showing database-wide statistics about query cancels occurring due to conflicts with recovery on standby servers. This view will only contain information on standby servers, since conflicts do not occur on master servers.

### Table 28-12. `pg_stat_all_tables` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>confl_deadlock</td>
<td>bigint</td>
<td>Number of queries in this database that have been canceled due to deadlocks</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of a table</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema that this table is in</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td>Name of this table</td>
</tr>
<tr>
<td>seq_scan</td>
<td>bigint</td>
<td>Number of sequential scans initiated on this table</td>
</tr>
<tr>
<td>seq_tup_read</td>
<td>bigint</td>
<td>Number of live rows fetched by sequential scans</td>
</tr>
<tr>
<td>idx_scan</td>
<td>bigint</td>
<td>Number of index scans initiated on this table</td>
</tr>
<tr>
<td>idx_tup_fetch</td>
<td>bigint</td>
<td>Number of live rows fetched by index scans</td>
</tr>
<tr>
<td>n_tup_ins</td>
<td>bigint</td>
<td>Number of rows inserted</td>
</tr>
<tr>
<td>n_tup_upd</td>
<td>bigint</td>
<td>Number of rows updated (includes HOT updated rows)</td>
</tr>
<tr>
<td>n_tup_del</td>
<td>bigint</td>
<td>Number of rows HOT deleted</td>
</tr>
<tr>
<td>n_tup_hot_upd</td>
<td>bigint</td>
<td>Number of rows HOT updated (i.e., with no separate index update required)</td>
</tr>
<tr>
<td>n_live_tup</td>
<td>bigint</td>
<td>Estimated number of live rows</td>
</tr>
<tr>
<td>n_dead_tup</td>
<td>bigint</td>
<td>Estimated number of dead rows</td>
</tr>
<tr>
<td>n_mod_since_analyze</td>
<td>bigint</td>
<td>Estimated number of rows modified since this table was last analyzed</td>
</tr>
<tr>
<td>last_vacuum</td>
<td>timestamp with time zone</td>
<td>Last time at which this table was manually vacuumed (not counting VACUUM FULL)</td>
</tr>
<tr>
<td>last_autovacuum</td>
<td>timestamp with time zone</td>
<td>Last time at which this table was vacuumed by the autovacuum daemon</td>
</tr>
<tr>
<td>last_analyze</td>
<td>timestamp with time zone</td>
<td>Last time at which this table was manually analyzed</td>
</tr>
<tr>
<td>last_autoanalyze</td>
<td>timestamp with time zone</td>
<td>Last time at which this table was analyzed by the autovacuum daemon</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

The `pg_stat_all_tables` view will contain one row for each table in the current database (including TOAST tables), showing statistics about accesses to that specific table. The `pg_stat_user_tables` and `pg_stat_sys_tables` views contain the same information, but filtered to only show user and system tables respectively.

### Table 28-13. pg_stat_all_indexes View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of the table for this index</td>
</tr>
<tr>
<td>indexrelid</td>
<td>oid</td>
<td>OID of this index</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema this index is in</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td>Name of the table for this index</td>
</tr>
<tr>
<td>indexrelname</td>
<td>name</td>
<td>Name of this index</td>
</tr>
<tr>
<td>idx_scan</td>
<td>bigint</td>
<td>Number of index scans initiated on this index</td>
</tr>
<tr>
<td>idx_tup_read</td>
<td>bigint</td>
<td>Number of index entries returned by scans on this index</td>
</tr>
<tr>
<td>idx_tup_fetch</td>
<td>bigint</td>
<td>Number of live table rows fetched by simple index scans using this index</td>
</tr>
</tbody>
</table>

The `pg_stat_all_indexes` view will contain one row for each index in the current database, showing statistics about accesses to that specific index. The `pg_stat_user_indexes` and `pg_stat_sys_indexes` views contain the same information, but filtered to only show user and system indexes respectively.

Indexes can be used by simple index scans, “bitmap” index scans, and the optimizer. In a bitmap scan the output of several indexes can be combined via AND or OR rules, so it is difficult to associate individual heap row fetches with specific indexes when a bitmap scan is used. Therefore, a bitmap scan increments the `pg_stat_all_indexes.idx_tup_read` count(s) for the index(es) it uses, and it increments the `pg_stat_all_tables.idx_tup_fetch` count for the table, but it does not affect `pg_stat_all_indexes.idx_tup_fetch`. The optimizer also accesses indexes to check for supplied constants whose values are outside the recorded range of the optimizer statistics because the optimizer statistics might be stale.

**Note:** The `idx_tup_read` and `idx_tup_fetch` counts can be different even without any use
of bitmap scans, because `idx_tup_read` counts index entries retrieved from the index while `idx_tup_fetch` counts live rows fetched from the table. The latter will be less if any dead or not-yet-committed rows are fetched using the index, or if any heap fetches are avoided by means of an index-only scan.

### Table 28-14. `pg_statio_all_tables` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>relid</code></td>
<td><code>oid</code></td>
<td>OID of a table</td>
</tr>
<tr>
<td><code>schemaname</code></td>
<td><code>name</code></td>
<td>Name of the schema that this table is in</td>
</tr>
<tr>
<td><code>relname</code></td>
<td><code>name</code></td>
<td>Name of this table</td>
</tr>
<tr>
<td><code>heap_blks_read</code></td>
<td><code>bigint</code></td>
<td>Number of disk blocks read from this table</td>
</tr>
<tr>
<td><code>heap_blks_hit</code></td>
<td><code>bigint</code></td>
<td>Number of buffer hits in this table</td>
</tr>
<tr>
<td><code>idx_blks_read</code></td>
<td><code>bigint</code></td>
<td>Number of disk blocks read from all indexes on this table</td>
</tr>
<tr>
<td><code>idx_blks_hit</code></td>
<td><code>bigint</code></td>
<td>Number of buffer hits in all indexes on this table</td>
</tr>
<tr>
<td><code>toast_blks_read</code></td>
<td><code>bigint</code></td>
<td>Number of disk blocks read from this table's TOAST table (if any)</td>
</tr>
<tr>
<td><code>toast_blks_hit</code></td>
<td><code>bigint</code></td>
<td>Number of buffer hits in this table’s TOAST table (if any)</td>
</tr>
<tr>
<td><code>tidx_blks_read</code></td>
<td><code>bigint</code></td>
<td>Number of disk blocks read from this table’s TOAST table indexes (if any)</td>
</tr>
<tr>
<td><code>tidx_blks_hit</code></td>
<td><code>bigint</code></td>
<td>Number of buffer hits in this table’s TOAST table indexes (if any)</td>
</tr>
</tbody>
</table>

The `pg_statio_all_tables` view will contain one row for each table in the current database (including TOAST tables), showing statistics about I/O on that specific table. The `pg_statio_user_tables` and `pg_statio_sys_tables` views contain the same information, but filtered to only show user and system tables respectively.

### Table 28-15. `pg_statio_all_indexes` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>relid</code></td>
<td><code>oid</code></td>
<td>OID of the table for this index</td>
</tr>
<tr>
<td><code>indexrelid</code></td>
<td><code>oid</code></td>
<td>OID of this index</td>
</tr>
<tr>
<td><code>schemaname</code></td>
<td><code>name</code></td>
<td>Name of the schema this index is in</td>
</tr>
<tr>
<td><code>relname</code></td>
<td><code>name</code></td>
<td>Name of the table for this index</td>
</tr>
<tr>
<td><code>indexrelname</code></td>
<td><code>name</code></td>
<td>Name of this index</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>idx_blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from this index</td>
</tr>
<tr>
<td>idx_blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in this index</td>
</tr>
</tbody>
</table>

The `pg_statio_all_indexes` view will contain one row for each index in the current database, showing statistics about I/O on that specific index. The `pg_statio_user_indexes` and `pg_statio_sys_indexes` views contain the same information, but filtered to only show user and system indexes respectively.

Table 28-16. `pg_statio_all_sequences` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of a sequence</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema this sequence is in</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td>Name of this sequence</td>
</tr>
<tr>
<td>blks_read</td>
<td>bigint</td>
<td>Number of disk blocks read from this sequence</td>
</tr>
<tr>
<td>blks_hit</td>
<td>bigint</td>
<td>Number of buffer hits in this sequence</td>
</tr>
</tbody>
</table>

The `pg_statio_all_sequences` view will contain one row for each sequence in the current database, showing statistics about I/O on that specific sequence.

Table 28-17. `pg_stat_user_functions` View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>funcid</td>
<td>oid</td>
<td>OID of a function</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>Name of the schema this function is in</td>
</tr>
<tr>
<td>funcname</td>
<td>name</td>
<td>Name of this function</td>
</tr>
<tr>
<td>calls</td>
<td>bigint</td>
<td>Number of times this function has been called</td>
</tr>
<tr>
<td>total_time</td>
<td>double precision</td>
<td>Total time spent in this function and all other functions called by it, in milliseconds</td>
</tr>
<tr>
<td>self_time</td>
<td>double precision</td>
<td>Total time spent in this function itself, not including other functions called by it, in milliseconds</td>
</tr>
</tbody>
</table>

The `pg_stat_user_functions` view will contain one row for each tracked function, showing statistics about executions of that function. The track_functions parameter controls exactly which functions are tracked.
28.2.3. Statistics Functions

Other ways of looking at the statistics can be set up by writing queries that use the same underlying statistics access functions used by the standard views shown above. For details such as the functions’ names, consult the definitions of the standard views. (For example, in psql you could issue `\d+ pg_stat_activity`.) The access functions for per-database statistics take a database OID as an argument to identify which database to report on. The per-table and per-index functions take a table or index OID. The functions for per-function statistics take a function OID. Note that only tables, indexes, and functions in the current database can be seen with these functions.

Additional functions related to statistics collection are listed in Table 28-18.

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_backend_pid()</code></td>
<td>integer</td>
<td>Process ID of the server process handling the current session</td>
</tr>
<tr>
<td><code>pg_stat_get_activity(integer)</code></td>
<td>setof record</td>
<td>Returns a record of information about the backend with the specified PID, or one record for each active backend in the system if NULL is specified. The fields returned are a subset of those in the <code>pg_stat_activity</code> view.</td>
</tr>
<tr>
<td><code>pg_stat_get_snapshot_timestamp()</code></td>
<td>timestamp with time zone</td>
<td>Returns the timestamp of the current statistics snapshot</td>
</tr>
<tr>
<td><code>pg_stat_clear_snapshot()</code></td>
<td>void</td>
<td>Discard the current statistics snapshot</td>
</tr>
<tr>
<td><code>pg_stat_reset()</code></td>
<td>void</td>
<td>Reset all statistics counters for the current database to zero (requires superuser privileges by default, but EXECUTE for this function can be granted to others.)</td>
</tr>
<tr>
<td><code>pg_stat_reset_shared(text)</code></td>
<td>void</td>
<td>Reset some cluster-wide statistics counters to zero, depending on the argument (requires superuser privileges by default, but EXECUTE for this function can be granted to others). Calling <code>pg_stat_reset_shared('bgwriter')</code> will zero all the counters shown in the <code>pg_stat_bgwriter</code> view. Calling <code>pg_stat_reset_shared('archiver')</code> will zero all the counters shown in the <code>pg_stat_archiver</code> view.</td>
</tr>
</tbody>
</table>
Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_stat_reset_single_table_counters(oid)</td>
<td>void</td>
<td>Reset statistics for a single table or index in the current database to zero (requires superuser privileges by default, but EXECUTE for this function can be granted to others)</td>
</tr>
<tr>
<td>pg_stat_reset_single_function_counters(oid)</td>
<td>void</td>
<td>Reset statistics for a single function in the current database to zero (requires superuser privileges by default, but EXECUTE for this function can be granted to others)</td>
</tr>
</tbody>
</table>

pg_stat_get_activity, the underlying function of the pg_stat_activity view, returns a set of records containing all the available information about each backend process. Sometimes it may be more convenient to obtain just a subset of this information. In such cases, an older set of per-backend statistics access functions can be used; these are shown in Table 28-19. These access functions use a backend ID number, which ranges from one to the number of currently active backends. The function pg_stat_get_backend_idset provides a convenient way to generate one row for each active backend for invoking these functions. For example, to show the PIDs and current queries of all backends:

```
SELECT pg_stat_get_backend_pid(s.backendid) AS pid,
       pg_stat_get_backend_activity(s.backendid) AS query
FROM (SELECT pg_stat_get_backend_idset() AS backendid) AS s;
```

| Table 28-19. Per-Backend Statistics Functions |
|-----------------------------------------------|------------------------------------------------|
| Function                                      | Return Type     | Description                                                                                                                                |
| pg_stat_get_backend_idset()                   | setof integer   | Set of currently active backend ID numbers (from 1 to the number of active backends)                                                        |
| pg_stat_get_backend_activity(integer)         | text            | Text of this backend’s most recent query                                                                                                  |
| pg_stat_get_backend_activity_start(integer)   | timestamp with time zone | Time when the most recent query was started                                                                                               |
| pg_stat_get_backend_client_addr(integer)      | inet            | IP address of the client connected to this backend                                                                                         |
| pg_stat_get_backend_client_port(integer)      | integer         | TCP port number that the client is using for communication                                                                                  |
| pg_stat_get_backend_dbid(integer)             | oid             | OID of the database this backend is connected to                                                                                           |
| pg_stat_get_backend_pid(integer)              | integer         | Process ID of this backend                                                                                                                 |
| pg_stat_get_backend_start(integer)            | timestamp with time zone | Time when this process was started                                                                                                         |
### Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_stat_get_backend_userid(integer)</code></td>
<td>oid</td>
<td>OID of the user logged into this backend</td>
</tr>
<tr>
<td><code>pg_stat_get_backend_wait_event_type(integer)</code></td>
<td>text</td>
<td>Wait event type name if backend is currently waiting, otherwise NULL. See Table 28-4 for details.</td>
</tr>
<tr>
<td><code>pg_stat_get_backend_wait_event(text,integer)</code></td>
<td>text</td>
<td>Wait event name if backend is currently waiting, otherwise NULL. See Table 28-4 for details.</td>
</tr>
<tr>
<td><code>pg_stat_get_backend_xact_start(text,timestamp)</code></td>
<td>text,timestamp with timezone</td>
<td>Time when the current transaction was started</td>
</tr>
</tbody>
</table>

#### 28.3. Viewing Locks

Another useful tool for monitoring database activity is the `pg_locks` system table. It allows the database administrator to view information about the outstanding locks in the lock manager. For example, this capability can be used to:

- View all the locks currently outstanding, all the locks on relations in a particular database, all the locks on a particular relation, or all the locks held by a particular PostgreSQL session.
- Determine the relation in the current database with the most ungranted locks (which might be a source of contention among database clients).
- Determine the effect of lock contention on overall database performance, as well as the extent to which contention varies with overall database traffic.

Details of the `pg_locks` view appear in Section 50.65. For more information on locking and managing concurrency with PostgreSQL, refer to Chapter 13.

#### 28.4. Progress Reporting

PostgreSQL has the ability to report the progress of certain commands during command execution. Currently, the only command which supports progress reporting is `VACUUM`. This may be expanded in the future.

##### 28.4.1. VACUUM Progress Reporting

Whenever `VACUUM` is running, the `pg_stat_progress_vacuum` view will contain one row for each backend (including autovacuum worker processes) that is currently vacuuming. The tables below describe the information that will be reported and provide information about how to interpret it. Progress reporting is not currently supported for `VACUUM FULL` and backends running `VACUUM FULL` will not be listed in this view.
### Table 28-20. pg_stat_progress_vacuum View

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pid</td>
<td>integer</td>
<td>Process ID of backend.</td>
</tr>
<tr>
<td>datid</td>
<td>oid</td>
<td>OID of the database to which this backend is connected.</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td>Name of the database to which this backend is connected.</td>
</tr>
<tr>
<td>relid</td>
<td>oid</td>
<td>OID of the table being vacuumed.</td>
</tr>
<tr>
<td>phase</td>
<td>text</td>
<td>Current processing phase of vacuum. See Table 28-21.</td>
</tr>
<tr>
<td>heap_blks_total</td>
<td>bigint</td>
<td>Total number of heap blocks in the table. This number is reported as of the beginning of the scan; blocks added later will not be (and need not be) visited by this VACUUM.</td>
</tr>
<tr>
<td>heap_blks_scanned</td>
<td>bigint</td>
<td>Number of heap blocks scanned. Because the visibility map is used to optimize scans, some blocks will be skipped without inspection; skipped blocks are included in this total, so that this number will eventually become equal to heap_blks_total when the vacuum is complete. This counter only advances when the phase is scanning heap.</td>
</tr>
<tr>
<td>heap_blks_vacuumed</td>
<td>bigint</td>
<td>Number of heap blocks vacuumed. Unless the table has no indexes, this counter only advances when the phase is vacuuming heap. Blocks that contain no dead tuples are skipped, so the counter may sometimes skip forward in large increments.</td>
</tr>
<tr>
<td>index_vacuum_count</td>
<td>bigint</td>
<td>Number of completed index vacuum cycles.</td>
</tr>
<tr>
<td>max_dead_tuples</td>
<td>bigint</td>
<td>Number of dead tuples that we can store before needing to perform an index vacuum cycle, based on maintenance_work_mem.</td>
</tr>
<tr>
<td>num_dead_tuples</td>
<td>bigint</td>
<td>Number of dead tuples collected since the last index vacuum cycle.</td>
</tr>
</tbody>
</table>
### Table 28-21. VACUUM phases

<table>
<thead>
<tr>
<th>Phase</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>initializing</td>
<td><strong>VACUUM</strong> is preparing to begin scanning the heap. This phase is expected to be very brief.</td>
</tr>
<tr>
<td>scanning heap</td>
<td><strong>VACUUM</strong> is currently scanning the heap. It will prune and defragment each page if required, and possibly perform freezing activity. The <code>heap_blks_scanned</code> column can be used to monitor the progress of the scan.</td>
</tr>
<tr>
<td>vacuuming indexes</td>
<td><strong>VACUUM</strong> is currently vacuuming the indexes. If a table has any indexes, this will happen at least once per vacuum, after the heap has been completely scanned. It may happen multiple times per vacuum if <code>maintenance_work_mem</code> is insufficient to store the number of dead tuples found.</td>
</tr>
<tr>
<td>vacuuming heap</td>
<td><strong>VACUUM</strong> is currently vacuuming the heap. Vacuuming the heap is distinct from scanning the heap, and occurs after each instance of vacuuming indexes. If <code>heap_blks_scanned</code> is less than <code>heap_blks_total</code>, the system will return to scanning the heap after this phase is completed; otherwise, it will begin cleaning up indexes after this phase is completed.</td>
</tr>
<tr>
<td>cleaning up indexes</td>
<td><strong>VACUUM</strong> is currently cleaning up indexes. This occurs after the heap has been completely scanned and all vacuuming of the indexes and the heap has been completed.</td>
</tr>
<tr>
<td>truncating heap</td>
<td><strong>VACUUM</strong> is currently truncating the heap so as to return empty pages at the end of the relation to the operating system. This occurs after cleaning up indexes.</td>
</tr>
<tr>
<td>performing final cleanup</td>
<td><strong>VACUUM</strong> is performing final cleanup. During this phase, <strong>VACUUM</strong> will vacuum the free space map, update statistics in <code>pg_class</code>, and report statistics to the statistics collector. When this phase is completed, <strong>VACUUM</strong> will end.</td>
</tr>
</tbody>
</table>

### 28.5. Dynamic Tracing

PostgreSQL provides facilities to support dynamic tracing of the database server. This allows an external utility to be called at specific points in the code and thereby trace execution.

A number of probes or trace points are already inserted into the source code. These probes are intended to be used by database developers and administrators. By default the probes are not compiled into PostgreSQL; the user needs to explicitly tell the configure script to make the probes available.
Currently, the DTrace\textsuperscript{1} utility is supported, which, at the time of this writing, is available on Solaris, OS X, FreeBSD, NetBSD, and Oracle Linux. The SystemTap\textsuperscript{2} project for Linux provides a DTrace equivalent and can also be used. Supporting other dynamic tracing utilities is theoretically possible by changing the definitions for the macros in \texttt{src/include/utils/probes.h}.

### 28.5.1. Compiling for Dynamic Tracing

By default, probes are not available, so you will need to explicitly tell the configure script to make the probes available in PostgreSQL. To include DTrace support specify \texttt{--enable-dtrace} to configure. See Section 16.4 for further information.

### 28.5.2. Built-in Probes

A number of standard probes are provided in the source code, as shown in Table 28-22; Table 28-23 shows the types used in the probes. More probes can certainly be added to enhance PostgreSQL’s observability.

**Table 28-22. Built-in DTrace Probes**

<table>
<thead>
<tr>
<th>Name</th>
<th>Parameters</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>transaction-start</td>
<td>(LocalTransactionId)</td>
<td>Probe that fires at the start of a new transaction. arg0 is the transaction ID.</td>
</tr>
<tr>
<td>transaction-commit</td>
<td>(LocalTransactionId)</td>
<td>Probe that fires when a transaction completes successfully. arg0 is the transaction ID.</td>
</tr>
<tr>
<td>transaction-abort</td>
<td>(LocalTransactionId)</td>
<td>Probe that fires when a transaction completes unsuccessfully. arg0 is the transaction ID.</td>
</tr>
<tr>
<td>query-start</td>
<td>(const char *)</td>
<td>Probe that fires when the processing of a query is started. arg0 is the query string.</td>
</tr>
<tr>
<td>query-done</td>
<td>(const char *)</td>
<td>Probe that fires when the processing of a query is complete. arg0 is the query string.</td>
</tr>
<tr>
<td>query-parse-start</td>
<td>(const char *)</td>
<td>Probe that fires when the parsing of a query is started. arg0 is the query string.</td>
</tr>
<tr>
<td>query-parse-done</td>
<td>(const char *)</td>
<td>Probe that fires when the parsing of a query is complete. arg0 is the query string.</td>
</tr>
</tbody>
</table>

1. \url{https://en.wikipedia.org/wiki/DTrace}
2. \url{http://sourceware.org/systemtap/}
<table>
<thead>
<tr>
<th>Name</th>
<th>Parameters</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>query-rewrite-start</td>
<td>(const char *)</td>
<td>Probe that fires when the rewriting of a query is started. arg0 is the query string.</td>
</tr>
<tr>
<td>query-rewrite-done</td>
<td>(const char *)</td>
<td>Probe that fires when the rewriting of a query is complete. arg0 is the query string.</td>
</tr>
<tr>
<td>query-plan-start</td>
<td>()</td>
<td>Probe that fires when the planning of a query is started.</td>
</tr>
<tr>
<td>query-plan-done</td>
<td>()</td>
<td>Probe that fires when the planning of a query is complete.</td>
</tr>
<tr>
<td>query-execute-start</td>
<td>()</td>
<td>Probe that fires when the execution of a query is started.</td>
</tr>
<tr>
<td>query-execute-done</td>
<td>()</td>
<td>Probe that fires when the execution of a query is complete.</td>
</tr>
<tr>
<td>statement-status</td>
<td>(const char *)</td>
<td>Probe that fires anytime the server process updates its pg_stat_activity.status. arg0 is the new status string.</td>
</tr>
<tr>
<td>checkpoint-start</td>
<td>(int)</td>
<td>Probe that fires when a checkpoint is started. arg0 holds the bitwise flags used to distinguish different checkpoint types, such as shutdown, immediate or force.</td>
</tr>
<tr>
<td>checkpoint-done</td>
<td>(int, int, int, int, int)</td>
<td>Probe that fires when a checkpoint is complete. (The probes listed next fire in sequence during checkpoint processing.) arg0 is the number of buffers written. arg1 is the total number of buffers. arg2, arg3 and arg4 contain the number of WAL files added, removed and recycled respectively.</td>
</tr>
<tr>
<td>clog-checkpoint-start</td>
<td>(bool)</td>
<td>Probe that fires when the CLOG portion of a checkpoint is started. arg0 is true for normal checkpoint, false for shutdown checkpoint.</td>
</tr>
<tr>
<td>clog-checkpoint-done</td>
<td>(bool)</td>
<td>Probe that fires when the CLOG portion of a checkpoint is complete. arg0 has the same meaning as for clog-checkpoint-start.</td>
</tr>
</tbody>
</table>
### Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Name</th>
<th>Parameters</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>subtrans-checkpoint-start</td>
<td>(bool)</td>
<td>Probe that fires when the SUBTRANS portion of a checkpoint is started. arg0 is true for normal checkpoint, false for shutdown checkpoint.</td>
</tr>
<tr>
<td>subtrans-checkpoint-done</td>
<td>(bool)</td>
<td>Probe that fires when the SUBTRANS portion of a checkpoint is complete. arg0 has the same meaning as for subtrans-checkpoint-start.</td>
</tr>
<tr>
<td>multixact-checkpoint-start</td>
<td>(bool)</td>
<td>Probe that fires when the MultiXact portion of a checkpoint is started. arg0 is true for normal checkpoint, false for shutdown checkpoint.</td>
</tr>
<tr>
<td>multixact-checkpoint-done</td>
<td>(bool)</td>
<td>Probe that fires when the MultiXact portion of a checkpoint is complete. arg0 has the same meaning as for multixact-checkpoint-start.</td>
</tr>
<tr>
<td>buffer-checkpoint-start</td>
<td>(int)</td>
<td>Probe that fires when the buffer-writing portion of a checkpoint is started. arg0 holds the bitwise flags used to distinguish different checkpoint types, such as shutdown, immediate or force.</td>
</tr>
<tr>
<td>buffer-sync-start</td>
<td>(int, int)</td>
<td>Probe that fires when we begin to write dirty buffers during checkpoint (after identifying which buffers must be written). arg0 is the total number of buffers. arg1 is the number that are currently dirty and need to be written.</td>
</tr>
<tr>
<td>buffer-sync-written</td>
<td>(int)</td>
<td>Probe that fires after each buffer is written during checkpoint. arg0 is the ID number of the buffer.</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------</td>
<td>-----------------------------------</td>
<td>-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>buffer-sync-done</td>
<td>(int, int, int)</td>
<td>Probe that fires when all dirty buffers have been written. arg0 is the total number of buffers. arg1 is the number of buffers actually written by the checkpoint process. arg2 is the number that were expected to be written (arg1 of buffer-sync-start); any difference reflects other processes flushing buffers during the checkpoint.</td>
</tr>
<tr>
<td>buffer-checkpoint-sync-start</td>
<td></td>
<td>Probe that fires after dirty buffers have been written to the kernel, and before starting to issue fsync requests.</td>
</tr>
<tr>
<td>buffer-checkpoint-done</td>
<td>()</td>
<td>Probe that fires when syncing of buffers to disk is complete.</td>
</tr>
<tr>
<td>twophase-checkpoint-start</td>
<td>()</td>
<td>Probe that fires when the two-phase portion of a checkpoint is started.</td>
</tr>
<tr>
<td>twophase-checkpoint-done</td>
<td>()</td>
<td>Probe that fires when the two-phase portion of a checkpoint is complete.</td>
</tr>
<tr>
<td>buffer-read-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, bool)</td>
<td>Probe that fires when a buffer read is started. arg0 and arg1 contain the fork and block numbers of the page (but arg1 will be -1 if this is a relation extension request). arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer. arg6 is true for a relation extension request, false for normal read.</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>---------------------------------------------------------------------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>buffer-read-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, bool, bool)</td>
<td>Probe that fires when a buffer read is complete. arg0 and arg1 contain the fork and block numbers of the page (if this is a relation extension request, arg1 now contains the block number of the newly added block). arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or <code>InvalidBackendId</code> (-1) for a shared buffer. arg6 is true for a relation extension request, false for normal read. arg7 is true if the buffer was found in the pool, false if not.</td>
</tr>
<tr>
<td>buffer-flush-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires before issuing any write request for a shared buffer. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation.</td>
</tr>
<tr>
<td>buffer-flush-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires when a write request is complete. (Note that this just reflects the time to pass the data to the kernel; it’s typically not actually been written to disk yet.) The arguments are the same as for <code>buffer-flush-start</code>.</td>
</tr>
<tr>
<td>buffer-write-dirty-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires when a server process begins to write a dirty buffer. (If this happens often, it implies that <code>shared_buffers</code> is too small or the background writer control parameters need adjustment.) arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation.</td>
</tr>
</tbody>
</table>
### Chapter 28. Monitoring Database Activity

<table>
<thead>
<tr>
<th>Name</th>
<th>Parameters</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>buffer-write-dirty-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid)</td>
<td>Probe that fires when a dirty-buffer write is complete. The arguments are the same as for buffer-write-dirty-start.</td>
</tr>
<tr>
<td>wal-buffer-write-dirty-start</td>
<td></td>
<td>Probe that fires when a server process begins to write a dirty WAL buffer because no more WAL buffer space is available. (If this happens often, it implies that wal_buffers is too small.)</td>
</tr>
<tr>
<td>wal-buffer-write-dirty-done</td>
<td></td>
<td>Probe that fires when a dirty WAL buffer write is complete.</td>
</tr>
<tr>
<td>xlog-insert</td>
<td>(unsigned char, unsigned char)</td>
<td>Probe that fires when a WAL record is inserted. arg0 is the resource manager (rmid) for the record. arg1 contains the info flags.</td>
</tr>
<tr>
<td>xlog-switch</td>
<td>()</td>
<td>Probe that fires when a WAL segment switch is requested.</td>
</tr>
<tr>
<td>smgr-md-read-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int)</td>
<td>Probe that fires when beginning to read a block from a relation. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer.</td>
</tr>
<tr>
<td>smgr-md-read-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, int, int)</td>
<td>Probe that fires when a block read is complete. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer. arg6 is the number of bytes actually read, while arg7 is the number requested (if these are different it indicates trouble).</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------</td>
<td>----------------------------------------------------------------------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>smgr-md-write-start</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int)</td>
<td>Probe that fires when beginning to write a block to a relation. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer.</td>
</tr>
<tr>
<td>smgr-md-write-done</td>
<td>(ForkNumber, BlockNumber, Oid, Oid, Oid, int, int, int)</td>
<td>Probe that fires when a block write is complete. arg0 and arg1 contain the fork and block numbers of the page. arg2, arg3, and arg4 contain the tablespace, database, and relation OIDs identifying the relation. arg5 is the ID of the backend which created the temporary relation for a local buffer, or InvalidBackendId (-1) for a shared buffer. arg6 is the number of bytes actually written, while arg7 is the number requested (if these are different it indicates trouble).</td>
</tr>
<tr>
<td>sort-start</td>
<td>(int, bool, int, int, bool)</td>
<td>Probe that fires when a sort operation is started. arg0 indicates heap, index or datum sort. arg1 is true for unique-value enforcement. arg2 is the number of key columns. arg3 is the number of kilobytes of work memory allowed. arg4 is true if random access to the sort result is required.</td>
</tr>
<tr>
<td>sort-done</td>
<td>(bool, long)</td>
<td>Probe that fires when a sort is complete. arg0 is true for external sort, false for internal sort. arg1 is the number of disk blocks used for an external sort, or kilobytes of memory used for an internal sort.</td>
</tr>
<tr>
<td>Name</td>
<td>Parameters</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------</td>
<td>-------------------------------------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>lwlock-acquire</td>
<td>(char *, int, LWLockMode)</td>
<td>Probe that fires when an LWLock has been acquired. arg0 is the LWLock’s tranche. arg1 is the LWLock’s offset within its tranche. arg2 is the requested lock mode, either exclusive or shared.</td>
</tr>
<tr>
<td>lwlock-release</td>
<td>(char *, int)</td>
<td>Probe that fires when an LWLock has been released (but note that any released waiters have not yet been awakened). arg0 is the LWLock’s tranche. arg1 is the LWLock’s offset within its tranche.</td>
</tr>
<tr>
<td>lwlock-wait-start</td>
<td>(char *, int, LWLockMode)</td>
<td>Probe that fires when an LWLock was not immediately available and a server process has begun to wait for the lock to become available. arg0 is the LWLock’s tranche. arg1 is the LWLock’s offset within its tranche. arg2 is the requested lock mode, either exclusive or shared.</td>
</tr>
<tr>
<td>lwlock-wait-done</td>
<td>(char *, int, LWLockMode)</td>
<td>Probe that fires when a server process has been released from its wait for an LWLock (it does not actually have the lock yet). arg0 is the LWLock’s tranche. arg1 is the LWLock’s offset within its tranche. arg2 is the requested lock mode, either exclusive or shared.</td>
</tr>
<tr>
<td>lwlock-condacquire</td>
<td>(char *, int, LWLockMode)</td>
<td>Probe that fires when an LWLock was successfully acquired when the caller specified no waiting. arg0 is the LWLock’s tranche. arg1 is the LWLock’s offset within its tranche. arg2 is the requested lock mode, either exclusive or shared.</td>
</tr>
</tbody>
</table>
Table 28-23. Defined Types Used in Probe Parameters

<table>
<thead>
<tr>
<th>Type</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>LocalTransactionId</td>
<td>unsigned int</td>
</tr>
<tr>
<td>LWLockMode</td>
<td>int</td>
</tr>
<tr>
<td>LOCKMODE</td>
<td>int</td>
</tr>
<tr>
<td>BlockNumber</td>
<td>unsigned int</td>
</tr>
<tr>
<td>Oid</td>
<td>unsigned int</td>
</tr>
<tr>
<td>ForkNumber</td>
<td>int</td>
</tr>
<tr>
<td>bool</td>
<td>char</td>
</tr>
</tbody>
</table>

28.5.3. Using Probes

The example below shows a DTrace script for analyzing transaction counts in the system, as an alternative to snapshotting pg_stat_database before and after a performance test:

```
#!/usr/sbin/dtrace -qs
```

Probe that fires when an LWLock was not successfully acquired when the caller specified no waiting. arg0 is the LWLock's tranche. arg1 is the LWLock's offset within its tranche. arg2 is the requested lock mode, either exclusive or shared.

Probe that fires when a request for a heavyweight lock (lmgr lock) has begun to wait because the lock is not available. arg0 through arg3 are the tag fields identifying the object being locked. arg4 indicates the type of object being locked. arg5 indicates the lock type being requested.

Probe that fires when a request for a heavyweight lock (lmgr lock) has finished waiting (i.e., has acquired the lock). The arguments are the same as for lock-wait-start.

Probe that fires when a deadlock is found by the deadlock detector.
Chapter 28. Monitoring Database Activity

When executed, the example D script gives output such as:

```
# ./txn_count.d 'pgrep -n postgres' or ./txn_count.d <PID>
^C
Start 71
Commit 70
Total time (ns) 2312105013
```

**Note:** SystemTap uses a different notation for trace scripts than DTrace does, even though the underlying trace points are compatible. One point worth noting is that at this writing, SystemTap scripts must reference probe names using double underscores in place of hyphens. This is expected to be fixed in future SystemTap releases.

You should remember that DTrace scripts need to be carefully written and debugged, otherwise the trace information collected might be meaningless. In most cases where problems are found it is the instrumentation that is at fault, not the underlying system. When discussing information found using dynamic tracing, be sure to enclose the script used to allow that too to be checked and discussed.

### 28.5.4. Defining New Probes

New probes can be defined within the code wherever the developer desires, though this will require a recompilation. Below are the steps for inserting new probes:

1. Decide on probe names and data to be made available through the probes
2. Add the probe definitions to `src/backend/utils/probes.d`
3. Include `pg_trace.h` if it is not already present in the module(s) containing the probe points, and insert `TRACE_POSTGRESQL` probe macros at the desired locations in the source code
4. Recompile and verify that the new probes are available
Example: Here is an example of how you would add a probe to trace all new transactions by transaction ID.

1. Decide that the probe will be named `transaction-start` and requires a parameter of type `LocalTransactionId`.

2. Add the probe definition to `src/backend/utils/probes.d`:
   ```c
   probe transaction_start(LocalTransactionId);
   ```
   Note the use of the double underline in the probe name. In a DTrace script using the probe, the double underline needs to be replaced with a hyphen, so `transaction-start` is the name to document for users.

3. At compile time, `transaction_start` is converted to a macro called `TRACE_POSTGRESQL_TRANSACTION_START` (notice the underscores are single here), which is available by including `pg_trace.h`. Add the macro call to the appropriate location in the source code. In this case, it looks like the following:
   ```c
   TRACE_POSTGRESQL_TRANSACTION_START(vxid.localTransactionId);
   ```

4. After recompiling and running the new binary, check that your newly added probe is available by executing the following DTrace command. You should see similar output:
   ```
   # dtrace -ln transaction-start
   ID PROVIDER MODULE FUNCTION NAME
   18705 postgresql49878 postgres StartTransactionCommand transaction-start
   18755 postgresql49877 postgres StartTransactionCommand transaction-start
   18805 postgresql49876 postgres StartTransactionCommand transaction-start
   18855 postgresql49875 postgres StartTransactionCommand transaction-start
   18986 postgresql49873 postgres StartTransactionCommand transaction-start
   ```

There are a few things to be careful about when adding trace macros to the C code:

- You should take care that the data types specified for a probe’s parameters match the data types of the variables used in the macro. Otherwise, you will get compilation errors.

- On most platforms, if PostgreSQL is built with `--enable-dtrace`, the arguments to a trace macro will be evaluated whenever control passes through the macro, *even if no tracing is being done*. This is usually not worth worrying about if you are just reporting the values of a few local variables. But beware of putting expensive function calls into the arguments. If you need to do that, consider protecting the macro with a check to see if the trace is actually enabled:
  ```c
  if (TRACE_POSTGRESQL_TRANSACTION_START_ENABLED())
      TRACE_POSTGRESQL_TRANSACTION_START(some_function(...));
  ```
  Each trace macro has a corresponding `ENABLED` macro.
Chapter 29. Monitoring Disk Usage

This chapter discusses how to monitor the disk usage of a PostgreSQL database system.

29.1. Determining Disk Usage

Each table has a primary heap disk file where most of the data is stored. If the table has any columns with potentially-wide values, there also might be a TOAST file associated with the table, which is used to store values too wide to fit comfortably in the main table (see Section 65.2). There will be one valid index on the TOAST table, if present. There also might be indexes associated with the base table. Each table and index is stored in a separate disk file — possibly more than one file, if the file would exceed one gigabyte. Naming conventions for these files are described in Section 65.1.

You can monitor disk space in three ways: using the SQL functions listed in Table 9-83, using the oid2name module, or using manual inspection of the system catalogs. The SQL functions are the easiest to use and are generally recommended. The remainder of this section shows how to do it by inspection of the system catalogs.

Using psql on a recently vacuumed or analyzed database, you can issue queries to see the disk usage of any table:

```sql
SELECT pg_relation_filepath(oid), relpages FROM pg_class WHERE relname = 'customer';
```

```
pg_relation_filepath | relpages
---------------------+----------
base/16384/16806     | 60
(1 row)
```

Each page is typically 8 kilobytes. (Remember, relpages is only updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE INDEX.) The file path name is of interest if you want to examine the table’s disk file directly.

To show the space used by TOAST tables, use a query like the following:

```sql
SELECT relname, relpages
FROM pg_class,
     (SELECT reltoastrelid
      FROM pg_class
      WHERE relname = 'customer') AS ss
WHERE oid = ss.reltoastrelid OR
      oid = (SELECT indexrelid
             FROM pg_index
             WHERE indrelid = ss.reltoastrelid)
ORDER BY relname;
```

```
relname | relpages
---------|--------
pg_toast_16806 | 0
pg_toast_16806_index | 1
```

You can easily display index sizes, too:

```sql
SELECT c2.relname, c2.relpages
FROM pg_class c, pg_class c2, pg_index i
```
WHERE c.relname = 'customer' AND
c.oid = i.indrelid AND
c2.oid = i.indexrelid
ORDER BY c2.relname;

<table>
<thead>
<tr>
<th>relname</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>customer_id_indexdex</td>
<td>26</td>
</tr>
</tbody>
</table>

It is easy to find your largest tables and indexes using this information:

SELECT relname, relpages
FROM pg_class
ORDER BY relpages DESC;

<table>
<thead>
<tr>
<th>relname</th>
<th>relpages</th>
</tr>
</thead>
<tbody>
<tr>
<td>bigtable</td>
<td>3290</td>
</tr>
<tr>
<td>customer</td>
<td>3144</td>
</tr>
</tbody>
</table>

### 29.2. Disk Full Failure

The most important disk monitoring task of a database administrator is to make sure the disk doesn’t become full. A filled data disk will not result in data corruption, but it might prevent useful activity from occurring. If the disk holding the WAL files grows full, database server panic and consequent shutdown might occur.

If you cannot free up additional space on the disk by deleting other things, you can move some of the database files to other file systems by making use of tablespaces. See Section 22.6 for more information about that.

**Tip:** Some file systems perform badly when they are almost full, so do not wait until the disk is completely full to take action.

If your system supports per-user disk quotas, then the database will naturally be subject to whatever quota is placed on the user the server runs as. Exceeding the quota will have the same bad effects as running out of disk space entirely.
Chapter 30. Reliability and the Write-Ahead Log

This chapter explains how the Write-Ahead Log is used to obtain efficient, reliable operation.

30.1. Reliability

Reliability is an important property of any serious database system, and PostgreSQL does everything possible to guarantee reliable operation. One aspect of reliable operation is that all data recorded by a committed transaction should be stored in a nonvolatile area that is safe from power loss, operating system failure, and hardware failure (except failure of the nonvolatile area itself, of course). Successfully writing the data to the computer’s permanent storage (disk drive or equivalent) ordinarily meets this requirement. In fact, even if a computer is fatally damaged, if the disk drives survive they can be moved to another computer with similar hardware and all committed transactions will remain intact.

While forcing data to the disk platters periodically might seem like a simple operation, it is not. Because disk drives are dramatically slower than main memory and CPUs, several layers of caching exist between the computer’s main memory and the disk platters. First, there is the operating system’s buffer cache, which caches frequently requested disk blocks and combines disk writes. Fortunately, all operating systems give applications a way to force writes from the buffer cache to disk, and PostgreSQL uses those features. (See the wal_sync_method parameter to adjust how this is done.)

Next, there might be a cache in the disk drive controller cards. Some of these caches are write-through, meaning writes are sent to the drive as soon as they arrive. Others are write-back, meaning data is sent to the drive at some later time. Such caches can be a reliability hazard because the memory in the disk controller cache is volatile, and will lose its contents in a power failure. Better controller cards have battery-backup units (BBUs), meaning the card has a battery that maintains power to the cache in case of system power loss. After power is restored the data will be written to the disk drives.

And finally, most disk drives have caches. Some are write-through while some are write-back, and the same concerns about data loss exist for write-back drive caches as for disk controller caches. Consumer-grade IDE and SATA drives are particularly likely to have write-back caches that will not survive a power failure. Many solid-state drives (SSD) also have volatile write-back caches.

These caches can typically be disabled; however, the method for doing this varies by operating system and drive type:

- On Linux, IDE and SATA drives can be queried using hdparm -I; write caching is enabled if there is a * next to Write cache. hdparm -W 0 can be used to turn off write caching. SCSI drives can be queried using sdparm. Use sdparm --get=WCE to check whether the write cache is enabled and sdparm --clear=WCE to disable it.

- On FreeBSD, IDE drives can be queried using atacontrol and write caching turned off using hw.ata.wc=0 in /boot/loader.conf; SCSI drives can be queried using camcontrol identify, and the write cache both queried and changed using sdparm when available.

- On Solaris, the disk write cache is controlled by format -e. (The Solaris ZFS file system is safe with disk write-cache enabled because it issues its own disk cache flush commands.)

Chapter 30. Reliability and the Write-Ahead Log

- On Windows, if `wal_sync_method` is `open_datsasync` (the default), write caching can be disabled by unchecking `My Computer\Open\disk drive\Properties\Hardware\Properties\Policies\Enable write caching on the disk`. Alternatively, set `wal_sync_method` to `fsync` or `fsync_writethrough`, which prevent write caching.

- On OS X, write caching can be prevented by setting `wal_sync_method` to `fsync_writethrough`.

Recent SATA drives (those following ATAPI-6 or later) offer a drive cache flush command (`FLUSH CACHE EXT`), while SCSI drives have long supported a similar command `SYNCHRONIZE CACHE`. These commands are not directly accessible to PostgreSQL, but some file systems (e.g., ZFS, ext4) can use them to flush data to the platters on write-back-enabled drives. Unfortunately, such file systems behave suboptimally when combined with battery-backup unit (BBU) disk controllers. In such setups, the synchronize command forces all data from the controller cache to the disks, eliminating much of the benefit of the BBU. You can run the `pg_test_fsync` program to see if you are affected. If you are affected, the performance benefits of the BBU can be regained by turning off write barriers in the file system or reconfiguring the disk controller, if that is an option. If write barriers are turned off, make sure the battery remains functional; a faulty battery can potentially lead to data loss. Hopefully file system and disk controller designers will eventually address this suboptimal behavior.

When the operating system sends a write request to the storage hardware, there is little it can do to make sure the data has arrived at a truly non-volatile storage area. Rather, it is the administrator’s responsibility to make certain that all storage components ensure integrity for both data and file-system metadata. Avoid disk controllers that have non-battery-backed write caches. At the drive level, disable write-back caching if the drive cannot guarantee the data will be written before shutdown. If you use SSDs, be aware that many of these do not honor cache flush commands by default. You can test for reliable I/O subsystem behavior using `diskchecker.pl`.

Another risk of data loss is posed by the disk platter write operations themselves. Disk platters are divided into sectors, commonly 512 bytes each. Every physical read or write operation processes a whole sector. When a write request arrives at the drive, it might be for some multiple of 512 bytes (PostgreSQL typically writes 8192 bytes, or 16 sectors, at a time), and the process of writing could fail due to power loss at any time, meaning some of the 512-byte sectors were written while others were not. To guard against such failures, PostgreSQL periodically writes full page images to permanent WAL storage before modifying the actual page on disk. By doing this, during crash recovery PostgreSQL can restore partially-written pages from WAL. If you have file-system software that prevents partial page writes (e.g., ZFS), you can turn off this page imaging by turning off the `full_page_writes` parameter. Battery-Backed Unit (BBU) disk controllers do not prevent partial page writes unless they guarantee that data is written to the BBU as full (8kB) pages.

PostgreSQL also protects against some kinds of data corruption on storage devices that may occur because of hardware errors or media failure over time, such as reading/writing garbage data.

- Each individual record in a WAL file is protected by a CRC-32 (32-bit) check that allows us to tell if record contents are correct. The CRC value is set when we write each WAL record and checked during crash recovery, archive recovery and replication.

- Data pages are not currently checksummed by default, though full page images recorded in WAL records will be protected; see `initdb` for details about enabling data page checksums.

- Internal data structures such as `pg_clog`, `pg_subtrans`, `pg_multixact`, `pg_serial`, `pg_notify`, `pg_stat`, `pg_snapshots` are not directly checksummed, nor are pages protected by full page writes. However, where such data structures are persistent, WAL records are written

---

that allow recent changes to be accurately rebuilt at crash recovery and those WAL records are protected as discussed above.

- Individual state files in `pg_twophase` are protected by CRC-32.
- Temporary data files used in larger SQL queries for sorts, materializations and intermediate results are not currently checksummed, nor will WAL records be written for changes to those files.

PostgreSQL does not protect against correctable memory errors and it is assumed you will operate using RAM that uses industry standard Error Correcting Codes (ECC) or better protection.

### 30.2. Write-Ahead Logging (WAL)

Write-Ahead Logging (WAL) is a standard method for ensuring data integrity. A detailed description can be found in most (if not all) books about transaction processing. Briefly, WAL's central concept is that changes to data files (where tables and indexes reside) must be written only after those changes have been logged, that is, after log records describing the changes have been flushed to permanent storage. If we follow this procedure, we do not need to flush data pages to disk on every transaction commit, because we know that in the event of a crash we will be able to recover the database using the log: any changes that have not been applied to the data pages can be redone from the log records. (This is roll-forward recovery, also known as REDO.)

**Tip:** Because WAL restores database file contents after a crash, journaled file systems are not necessary for reliable storage of the data files or WAL files. In fact, journaling overhead can reduce performance, especially if journaling causes file system data to be flushed to disk. Fortunately, data flushing during journaling can often be disabled with a file system mount option, e.g. `data=writeback` on a Linux ext3 file system. Journaled file systems do improve boot speed after a crash.

Using WAL results in a significantly reduced number of disk writes, because only the log file needs to be flushed to disk to guarantee that a transaction is committed, rather than every data file changed by the transaction. The log file is written sequentially, and so the cost of syncing the log is much less than the cost of flushing the data pages. This is especially true for servers handling many small transactions touching different parts of the data store. Furthermore, when the server is processing many small concurrent transactions, one `fsync` of the log file may suffice to commit many transactions.

WAL also makes it possible to support on-line backup and point-in-time recovery, as described in Section 25.3. By archiving the WAL data we can support reverting to any time instant covered by the available WAL data: we simply install a prior physical backup of the database, and replay the WAL log just as far as the desired time. What’s more, the physical backup doesn’t have to be an instantaneous snapshot of the database state — if it is made over some period of time, then replaying the WAL log for that period will fix any internal inconsistencies.

### 30.3. Asynchronous Commit

Asynchronous commit is an option that allows transactions to complete more quickly, at the cost that the most recent transactions may be lost if the database should crash. In many applications this is an acceptable trade-off.
As described in the previous section, transaction commit is normally synchronous: the server waits for the transaction’s WAL records to be flushed to permanent storage before returning a success indication to the client. The client is therefore guaranteed that a transaction reported to be committed will be preserved, even in the event of a server crash immediately after. However, for short transactions this delay is a major component of the total transaction time. Selecting asynchronous commit mode means that the server returns success as soon as the transaction is logically completed, before the WAL records it generated have actually made their way to disk. This can provide a significant boost in throughput for small transactions.

Asynchronous commit introduces the risk of data loss. There is a short time window between the report of transaction completion to the client and the time that the transaction is truly committed (that is, it is guaranteed not to be lost if the server crashes). Thus asynchronous commit should not be used if the client will take external actions relying on the assumption that the transaction will be remembered. As an example, a bank would certainly not use asynchronous commit for a transaction recording an ATM’s dispensing of cash. But in many scenarios, such as event logging, there is no need for a strong guarantee of this kind.

The user can select the commit mode of each transaction, so that it is possible to have both synchronous and asynchronous commit transactions running concurrently. This allows flexible trade-offs between performance and certainty of transaction durability. The commit mode is controlled by the user-settable parameter synchronous_commit, which can be changed in any of the ways that a configuration parameter can be set. The mode used for any one transaction depends on the value of synchronous_commit when transaction commit begins.

Certain utility commands, for instance DROP TABLE, are forced to commit synchronously regardless of the setting of synchronous_commit. This is to ensure consistency between the server’s file system and the logical state of the database. The commands supporting two-phase commit, such as PREPARE TRANSACTION, are also always synchronous.

If the database crashes during the risk window between an asynchronous commit and the writing of the transaction’s WAL records, then changes made during that transaction will be lost. The duration of the risk window is limited because a background process (the “WAL writer”) flushes unwritten WAL records to disk every wal_writer_delay milliseconds. The actual maximum duration of the risk window is three times wal_writer_delay because the WAL writer is designed to favor writing whole pages at a time during busy periods.

Caution

An immediate-mode shutdown is equivalent to a server crash, and will therefore cause loss of any unflushed asynchronous commits.

Asynchronous commit provides behavior different from setting fsync = off. fsync is a server-wide setting that will alter the behavior of all transactions. It disables all logic within PostgreSQL that attempts to synchronize writes to different portions of the database, and therefore a system crash (that is, a hardware or operating system crash, not a failure of PostgreSQL itself) could result in arbitrarily bad corruption of the database state. In many scenarios, asynchronous commit provides most of the
performance improvement that could be obtained by turning off `fsync`, but without the risk of data corruption.

`commit_delay` also sounds very similar to asynchronous commit, but it is actually a synchronous commit method (in fact, `commit_delay` is ignored during an asynchronous commit). `commit_delay` causes a delay just before a transaction flushes WAL to disk, in the hope that a single flush executed by one such transaction can also serve other transactions committing at about the same time. The setting can be thought of as a way of increasing the time window in which transactions can join a group about to participate in a single flush, to amortize the cost of the flush among multiple transactions.

### 30.4. WAL Configuration

There are several WAL-related configuration parameters that affect database performance. This section explains their use. Consult Chapter 19 for general information about setting server configuration parameters.

**Checkpoints** are points in the sequence of transactions at which it is guaranteed that the heap and index data files have been updated with all information written before that checkpoint. At checkpoint time, all dirty data pages are flushed to disk and a special checkpoint record is written to the log file. (The change records were previously flushed to the WAL files.) In the event of a crash, the crash recovery procedure looks at the latest checkpoint record to determine the point in the log (known as the redo record) from which it should start the REDO operation. Any changes made to data files before that point are guaranteed to be already on disk. Hence, after a checkpoint, log segments preceding the one containing the redo record are no longer needed and can be recycled or removed. (When WAL archiving is being done, the log segments must be archived before being recycled or removed.)

The checkpoint requirement of flushing all dirty data pages to disk can cause a significant I/O load. For this reason, checkpoint activity is throttled so that I/O begins at checkpoint start and completes before the next checkpoint is due to start; this minimizes performance degradation during checkpoints.

The server’s checkpoint process automatically performs a checkpoint every so often. A checkpoint is begun every `checkpoint_timeout` seconds, or if `max_wal_size` is about to be exceeded, whichever comes first. The default settings are 5 minutes and 1 GB, respectively. If no WAL has been written since the previous checkpoint, new checkpoints will be skipped even if `checkpoint_timeout` has passed. (If WAL archiving is being used and you want to put a lower limit on how often files are archived in order to bound potential data loss, you should adjust the `archive_timeout` parameter rather than the checkpoint parameters.) It is also possible to force a checkpoint by using the SQL command `CHECKPOINT`.

Reducing `checkpoint_timeout` and/or `max_wal_size` causes checkpoints to occur more often. This allows faster after-crash recovery, since less work will need to be redone. However, one must balance this against the increased cost of flushing dirty data pages more often. If `full_page_writes` is set (as is the default), there is another factor to consider. To ensure data page consistency, the first modification of a data page after each checkpoint results in logging the entire page content. In that case, a smaller checkpoint interval increases the volume of output to the WAL log, partially negating the goal of using a smaller interval, and in any case causing more disk I/O.

Checkpoints are fairly expensive, first because they require writing out all currently dirty buffers, and second because they result in extra subsequent WAL traffic as discussed above. It is therefore wise to set the checkpointing parameters high enough so that checkpoints don’t happen too often. As a simple sanity check on your checkpointing parameters, you can set the `checkpoint_warning` parameter. If checkpoints happen closer together than `checkpoint_warning` seconds, a message will be output to the server log recommending increasing `max_wal_size`. Occasional appearance of such a message is not cause for alarm, but if it appears often then the checkpoint control parameters
should be increased. Bulk operations such as large `COPY` transfers might cause a number of such warnings to appear if you have not set `max_wal_size` high enough.

To avoid flooding the I/O system with a burst of page writes, writing dirty buffers during a checkpoint is spread over a period of time. That period is controlled by `checkpoint_completion_target`, which is given as a fraction of the checkpoint interval. The I/O rate is adjusted so that the checkpoint finishes when the given fraction of `checkpoint_timeout` seconds have elapsed, or before `max_wal_size` is exceeded, whichever is sooner. With the default value of 0.5, PostgreSQL can be expected to complete each checkpoint in about half the time before the next checkpoint starts. On a system that’s very close to maximum I/O throughput during normal operation, you might want to increase `checkpoint_completion_target` to reduce the I/O load from checkpoints. The disadvantage of this is that prolonging checkpoints affects recovery time, because more WAL segments will need to be kept around for possible use in recovery. Although `checkpoint_completion_target` can be set as high as 1.0, it is best to keep it less than that (perhaps 0.9 at most) since checkpoints include some other activities besides writing dirty buffers. A setting of 1.0 is quite likely to result in checkpoints not being completed on time, which would result in performance loss due to unexpected variation in the number of WAL segments needed.

On Linux and POSIX platforms `checkpoint_flush_after` allows to force the OS that pages written by the checkpoint should be flushed to disk after a configurable number of bytes. Otherwise, these pages may be kept in the OS’s page cache, inducing a stall when `fsync` is issued at the end of a checkpoint. This setting will often help to reduce transaction latency, but it also can have an adverse effect on performance; particularly for workloads that are bigger than `shared_buffers`, but smaller than the OS’s page cache.

The number of WAL segment files in `pg_xlog` directory depends on `min_wal_size`, `max_wal_size` and the amount of WAL generated in previous checkpoint cycles. When old log segment files are no longer needed, they are removed or recycled (that is, renamed to become future segments in the numbered sequence). If, due to a short-term peak of log output rate, `max_wal_size` is exceeded, the unneeded segment files will be removed until the system gets back under this limit. Below that limit, the system recycles enough WAL files to cover the estimated need until the next checkpoint, and removes the rest. The estimate is based on a moving average of the number of WAL files used in previous checkpoint cycles. The moving average is increased immediately if the actual usage exceeds the estimate, so it accommodates peak usage rather than average usage to some extent. `min_wal_size` puts a minimum on the amount of WAL files recycled for future use; that much WAL is always recycled for future use, even if the system is idle and the WAL usage estimate suggests that little WAL is needed.

Independently of `max_wal_size`, `wal_keep_segments + 1` most recent WAL files are kept at all times. Also, if WAL archiving is used, old segments can not be removed or recycled until they are archived. If WAL archiving cannot keep up with the pace that WAL is generated, or if `archive_command` fails repeatedly, old WAL files will accumulate in `pg_xlog` until the situation is resolved. A slow or failed standby server that uses a replication slot will have the same effect (see Section 26.2.6).

In archive recovery or standby mode, the server periodically performs `restartpoints`, which are similar to checkpoints in normal operation: the server forces all its state to disk, updates the `pg_control` file to indicate that the already-processed WAL data need not be scanned again, and then recycles any old log segment files in the `pg_xlog` directory. Restartpoints can’t be performed more frequently than checkpoints in the master because restartpoints can only be performed at checkpoint records. A restartpoint is triggered when a checkpoint record is reached if at least `checkpoint_timeout` seconds have passed since the last restartpoint, or if WAL size is about to exceed `max_wal_size`. However, because of limitations on when a restartpoint can be performed, `max_wal_size` is often exceeded during recovery, by up to one checkpoint cycle’s worth of WAL. (`max_wal_size` is never a hard limit anyway, so you should always leave plenty of headroom to avoid running out of disk space.)
Chapter 30. Reliability and the Write-Ahead Log

There are two commonly used internal WAL functions: XLogInsertRecord and XLogFlush. XLogInsertRecord is used to place a new record into the WAL buffers in shared memory. If there is no space for the new record, XLogInsertRecord will have to write (move to kernel cache) a few filled WAL buffers. This is undesirable because XLogInsertRecord is used on every database low level modification (for example, row insertion) at a time when an exclusive lock is held on affected data pages, so the operation needs to be as fast as possible. What is worse, writing WAL buffers might also force the creation of a new log segment, which takes even more time. Normally, WAL buffers should be written and flushed by an XLogFlush request, which is made, for the most part, at transaction commit time to ensure that transaction records are flushed to permanent storage. On systems with high log output, XLogFlush requests might not occur often enough to prevent XLogInsertRecord from having to do writes. On such systems one should increase the number of WAL buffers by modifying the wal_buffers parameter. When full_page_writes is set and the system is very busy, setting wal_buffers higher will help smooth response times during the period immediately following each checkpoint.

The commit_delay parameter defines for how many microseconds a group commit leader process will sleep after acquiring a lock within XLogFlush, while group commit followers queue up behind the leader. This delay allows other server processes to add their commit records to the WAL buffers so that all of them will be flushed by the leader’s eventual sync operation. No sleep will occur if fsync is not enabled, or if fewer than commit_siblings other sessions are currently in active transactions; this avoids sleeping when it’s unlikely that any other session will commit soon. Note that on some platforms, the resolution of a sleep request is ten milliseconds, so that any nonzero commit_delay setting between 1 and 10000 microseconds would have the same effect. Note also that on some platforms, sleep operations may take slightly longer than requested by the parameter.

Since the purpose of commit_delay is to allow the cost of each flush operation to be amortized across concurrently committing transactions (potentially at the expense of transaction latency), it is necessary to quantify that cost before the setting can be chosen intelligently. The higher that cost is, the more effective commit_delay is expected to be in increasing transaction throughput, up to a point. The pg_test_fsync program can be used to measure the average time in microseconds that a WAL flush operation takes. A value of half of the average time the program reports it takes to flush after a single 8kB write operation is often the most effective setting for commit_delay, so this value is recommended as the starting point to use when optimizing for a particular workload. While tuning commit_delay is particularly useful when the WAL log is stored on high-latency rotating disks, benefits can be significant even on storage media with very fast sync times, such as solid-state drives or RAID arrays with a battery-backed write cache; but this should definitely be tested against a representative workload. Higher values of commit_siblings should be used in such cases, whereas smaller commit_siblings values are often helpful on higher latency media. Note that it is quite possible that a setting of commit_delay that is too high can increase transaction latency by so much that total transaction throughput suffers.

When commit_delay is set to zero (the default), it is still possible for a form of group commit to occur, but each group will consist only of sessions that reach the point where they need to flush their commit records during the window in which the previous flush operation (if any) is occurring. At higher client counts a “gangway effect” tends to occur, so that the effects of group commit become significant even when commit_delay is zero, and thus explicitly setting commit_delay tends to help less. Setting commit_delay can only help when (1) there are some concurrently committing transactions, and (2) throughput is limited to some degree by commit rate; but with high rotational latency this setting can be effective in increasing transaction throughput with as few as two clients (that is, a single committing client with one sibling transaction).

The wal_sync_method parameter determines how PostgreSQL will ask the kernel to force WAL updates out to disk. All the options should be the same in terms of reliability, with the exception of fsync_writethrough, which can sometimes force a flush of the disk cache even when other op-
tions do not do so. However, it’s quite platform-specific which one will be the fastest. You can test the speeds of different options using the pg_test_fsync program. Note that this parameter is irrelevant if fsync has been turned off.

Enabling the wal_debug configuration parameter (provided that PostgreSQL has been compiled with support for it) will result in each XLogInsertRecord and XLogFlush WAL call being logged to the server log. This option might be replaced by a more general mechanism in the future.

### 30.5. WAL Internals

WAL is automatically enabled; no action is required from the administrator except ensuring that the disk-space requirements for the WAL logs are met, and that any necessary tuning is done (see Section 30.4).

WAL logs are stored in the directory `pg_xlog` under the data directory, as a set of segment files, normally each 16 MB in size (but the size can be changed by altering the `--with-wal-segsize` configure option when building the server). Each segment is divided into pages, normally 8 kB each (this size can be changed via the `--with-wal-blocksize` configure option). The log record headers are described in `access/xlogrecord.h`; the record content is dependent on the type of event that is being logged. Segment files are given ever-increasing numbers as names, starting at `000000010000000000000000`. The numbers do not wrap, but it will take a very, very long time to exhaust the available stock of numbers.

It is advantageous if the log is located on a different disk from the main database files. This can be achieved by moving the `pg_xlog` directory to another location (while the server is shut down, of course) and creating a symbolic link from the original location in the main data directory to the new location.

The aim of WAL is to ensure that the log is written before database records are altered, but this can be subverted by disk drives that falsely report a successful write to the kernel, when in fact they have only cached the data and not yet stored it on the disk. A power failure in such a situation might lead to irrecoverable data corruption. Administrators should try to ensure that disks holding PostgreSQL’s WAL log files do not make such false reports. (See Section 30.1.)

After a checkpoint has been made and the log flushed, the checkpoint’s position is saved in the file `pg_control`. Therefore, at the start of recovery, the server first reads `pg_control` and then the checkpoint record; then it performs the REDO operation by scanning forward from the log position indicated in the checkpoint record. Because the entire content of data pages is saved in the log on the first page modification after a checkpoint (assuming `full_page_writes` is not disabled), all pages changed since the checkpoint will be restored to a consistent state.

To deal with the case where `pg_control` is corrupt, we should support the possibility of scanning existing log segments in reverse order — newest to oldest — in order to find the latest checkpoint. This has not been implemented yet. `pg_control` is small enough (less than one disk page) that it is not subject to partial-write problems, and as of this writing there have been no reports of database failures due solely to the inability to read `pg_control` itself. So while it is theoretically a weak spot, `pg_control` does not seem to be a problem in practice.
Chapter 31. Regression Tests

The regression tests are a comprehensive set of tests for the SQL implementation in PostgreSQL. They test standard SQL operations as well as the extended capabilities of PostgreSQL.

31.1. Running the Tests

The regression tests can be run against an already installed and running server, or using a temporary installation within the build tree. Furthermore, there is a “parallel” and a “sequential” mode for running the tests. The sequential method runs each test script alone, while the parallel method starts up multiple server processes to run groups of tests in parallel. Parallel testing adds confidence that interprocess communication and locking are working correctly.

31.1.1. Running the Tests Against a Temporary Installation

To run the parallel regression tests after building but before installation, type:

```
make check
```

in the top-level directory. (Or you can change to `src/test/regress` and run the command there.) At the end you should see something like:

```
=======================
All 115 tests passed.
=======================
```

or otherwise a note about which tests failed. See Section 31.2 below before assuming that a “failure” represents a serious problem.

Because this test method runs a temporary server, it will not work if you did the build as the root user, since the server will not start as root. Recommended procedure is not to do the build as root, or else to perform testing after completing the installation.

If you have configured PostgreSQL to install into a location where an older PostgreSQL installation already exists, and you perform `make check` before installing the new version, you might find that the tests fail because the new programs try to use the already-installed shared libraries. (Typical symptoms are complaints about undefined symbols.) If you wish to run the tests before overwriting the old installation, you’ll need to build with `configure --disable-rpath`. It is not recommended that you use this option for the final installation, however.

The parallel regression test starts quite a few processes under your user ID. Presently, the maximum concurrency is twenty parallel test scripts, which means forty processes: there’s a server process and a psql process for each test script. So if your system enforces a per-user limit on the number of processes, make sure this limit is at least fifty or so, else you might get random-seeming failures in the parallel test. If you are not in a position to raise the limit, you can cut down the degree of parallelism by setting the `MAX_CONNECTIONS` parameter. For example:

```
make MAX_CONNECTIONS=10 check
```

runs no more than ten tests concurrently.
### 31.1.2. Running the Tests Against an Existing Installation

To run the tests after installation (see Chapter 16), initialize a data area and start the server as explained in Chapter 18, then type:

```
make installcheck
```

or for a parallel test:

```
make installcheck-parallel
```

The tests will expect to contact the server at the local host and the default port number, unless directed otherwise by `PGHOST` and `PGPORT` environment variables. The tests will be run in a database named `regression`; any existing database by this name will be dropped.

The tests will also transiently create some cluster-wide objects, such as roles and tablespaces. These objects will have names beginning with `regress_`. Beware of using `installcheck` mode in installations that have any actual users or tablespaces named that way.

### 31.1.3. Additional Test Suites

The `make check` and `make installcheck` commands run only the “core” regression tests, which test built-in functionality of the PostgreSQL server. The source distribution also contains additional test suites, most of them having to do with add-on functionality such as optional procedural languages.

To run all test suites applicable to the modules that have been selected to be built, including the core tests, type one of these commands at the top of the build tree:

```
make check-world
make installcheck-world
```

These commands run the tests using temporary servers or an already-installed server, respectively, just as previously explained for `make check` and `make installcheck`. Other considerations are the same as previously explained for each method. Note that `make check-world` builds a separate temporary installation tree for each tested module, so it requires a great deal more time and disk space than `make installcheck-world`.

Alternatively, you can run individual test suites by typing `make check` or `make installcheck` in the appropriate subdirectory of the build tree. Keep in mind that `make installcheck` assumes you’ve installed the relevant module(s), not only the core server.

The additional tests that can be invoked this way include:

- Regression tests for optional procedural languages (other than PL/pgSQL, which is tested by the core tests). These are located under `src/pl`.
- Regression tests for `contrib` modules, located under `contrib`. Not all `contrib` modules have tests.
- Regression tests for the ECPG interface library, located in `src/interfaces/ecpg/test`.
- Tests stressing behavior of concurrent sessions, located in `src/test/isolation`.
- Tests of client programs under `src/bin`. See also Section 31.4.

When using `installcheck` mode, these tests will destroy any existing databases named `pl_regression`, `contrib_regression`, `isolation_regression`, `ecpg1_regression`, or `ecpg2_regression`, as well as `regression`.  

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*719*
31.1.4. Locale and Encoding

By default, tests using a temporary installation use the locale defined in the current environment and the corresponding database encoding as determined by `initdb`. It can be useful to test different locales by setting the appropriate environment variables, for example:

```make
make check LANG=C
make check LC_COLLATE=en_US.utf8 LC_CTYPE=fr_CA.utf8
```

For implementation reasons, setting `LC_ALL` does not work for this purpose; all the other locale-related environment variables do work.

When testing against an existing installation, the locale is determined by the existing database cluster and cannot be set separately for the test run.

You can also choose the database encoding explicitly by setting the variable `ENCODING`, for example:

```make
make check LANG=C ENCODING=EUC_JP
```

Setting the database encoding this way typically only makes sense if the locale is C; otherwise the encoding is chosen automatically from the locale, and specifying an encoding that does not match the locale will result in an error.

The database encoding can be set for tests against either a temporary or an existing installation, though in the latter case it must be compatible with the installation’s locale.

31.1.5. Extra Tests

The core regression test suite contains a few test files that are not run by default, because they might be platform-dependent or take a very long time to run. You can run these or other extra test files by setting the variable `EXTRA_TESTS`. For example, to run the `numeric_big` test:

```make
make check EXTRA_TESTS=numeric_big
```

To run the collation tests:

```make
make check EXTRA_TESTS=collate.linux.utf8 LANG=en_US.utf8
```

The `collate.linux.utf8` test works only on Linux/glibc platforms, and only when run in a database that uses UTF-8 encoding.

31.1.6. Testing Hot Standby

The source distribution also contains regression tests for the static behavior of Hot Standby. These tests require a running primary server and a running standby server that is accepting new WAL changes from the primary (using either file-based log shipping or streaming replication). Those servers are not automatically created for you, nor is replication setup documented here. Please check the various sections of the documentation devoted to the required commands and related issues.

To run the Hot Standby tests, first create a database called `regression` on the primary:

```psql
psql -h primary -c "CREATE DATABASE regression"
```

Next, run the preparatory script `src/test/regress/sql/hs_primary_setup.sql` on the primary in the regression database, for example:
psql -h primary -f src/test/regress/sql/hs_primary_setup.sql regression

Allow these changes to propagate to the standby.

Now arrange for the default database connection to be to the standby server under test (for example, by setting the PGHOST and PGPORT environment variables). Finally, run `make standbycheck` in the regression directory:

cd src/test/regress
make standbycheck

Some extreme behaviors can also be generated on the primary using the script `src/test/regress/sql/hs_primary_extremes.sql` to allow the behavior of the standby to be tested.

31.2. Test Evaluation

Some properly installed and fully functional PostgreSQL installations can “fail” some of these regression tests due to platform-specific artifacts such as varying floating-point representation and message wording. The tests are currently evaluated using a simple `diff` comparison against the outputs generated on a reference system, so the results are sensitive to small system differences. When a test is reported as “failed”, always examine the differences between expected and actual results; you might find that the differences are not significant. Nonetheless, we still strive to maintain accurate reference files across all supported platforms, so it can be expected that all tests pass.

The actual outputs of the regression tests are in files in the `src/test/regress/results` directory. The test script uses `diff` to compare each output file against the reference outputs stored in the `src/test/regress/expected` directory. Any differences are saved for your inspection in `src/test/regress/regression.diffs`. (When running a test suite other than the core tests, these files of course appear in the relevant subdirectory, not `src/test/regress`.)

If you don’t like the `diff` options that are used by default, set the environment variable `PG_REGRESS_DIFF_OPTS`, for instance `PG_REGRESS_DIFF_OPTS=’-u’`. (Or you can run `diff` yourself, if you prefer.)

If for some reason a particular platform generates a “failure” for a given test, but inspection of the output convinces you that the result is valid, you can add a new comparison file to silence the failure report in future test runs. See Section 31.3 for details.

31.2.1. Error Message Differences

Some of the regression tests involve intentional invalid input values. Error messages can come from either the PostgreSQL code or from the host platform system routines. In the latter case, the messages can vary between platforms, but should reflect similar information. These differences in messages will result in a “failed” regression test that can be validated by inspection.

31.2.2. Locale Differences

If you run the tests against a server that was initialized with a collation-order locale other than C, then there might be differences due to sort order and subsequent failures. The regression test suite is set
up to handle this problem by providing alternate result files that together are known to handle a large number of locales.

To run the tests in a different locale when using the temporary-installation method, pass the appropriate locale-related environment variables on the `make` command line, for example:

```
make check LANG=de_DE.utf8
```

(The regression test driver unsets `LC_ALL`, so it does not work to choose the locale using that variable.) To use no locale, either unset all locale-related environment variables (or set them to `C`) or use the following special invocation:

```
make check NO_LOCALE=1
```

When running the tests against an existing installation, the locale setup is determined by the existing installation. To change it, initialize the database cluster with a different locale by passing the appropriate options to `initdb`.

In general, it is advisable to try to run the regression tests in the locale setup that is wanted for production use, as this will exercise the locale- and encoding-related code portions that will actually be used in production. Depending on the operating system environment, you might get failures, but then you will at least know what locale-specific behaviors to expect when running real applications.

### 31.2.3. Date and Time Differences

Most of the date and time results are dependent on the time zone environment. The reference files are generated for time zone `PST8PDT` (Berkeley, California), and there will be apparent failures if the tests are not run with that time zone setting. The regression test driver sets environment variable `PGTZ` to `PST8PDT`, which normally ensures proper results.

### 31.2.4. Floating-Point Differences

Some of the tests involve computing 64-bit floating-point numbers (double precision) from table columns. Differences in results involving mathematical functions of double precision columns have been observed. The `float8` and `geometry` tests are particularly prone to small differences across platforms, or even with different compiler optimization settings. Human eyeball comparison is needed to determine the real significance of these differences which are usually 10 places to the right of the decimal point.

Some systems display minus zero as `-0`, while others just show `0`.

Some systems signal errors from `pow()` and `exp()` differently from the mechanism expected by the current PostgreSQL code.

### 31.2.5. Row Ordering Differences

You might see differences in which the same rows are output in a different order than what appears in the expected file. In most cases this is not, strictly speaking, a bug. Most of the regression test scripts are not so pedantic as to use an `ORDER BY` for every single `SELECT`, and so their result row orderings are not well-defined according to the SQL specification. In practice, since we are looking at the same queries being executed on the same data by the same software, we usually get the same result ordering on all platforms, so the lack of `ORDER BY` is not a problem. Some queries do exhibit cross-platform
ordering differences, however. When testing against an already-installed server, ordering differences can also be caused by non-C locale settings or non-default parameter settings, such as custom values of `work_mem` or the planner cost parameters.

Therefore, if you see an ordering difference, it’s not something to worry about, unless the query does have an `ORDER BY` that your result is violating. However, please report it anyway, so that we can add an `ORDER BY` to that particular query to eliminate the bogus “failure” in future releases.

You might wonder why we don’t order all the regression test queries explicitly to get rid of this issue once and for all. The reason is that that would make the regression tests less useful, not more, since they’d tend to exercise query plan types that produce ordered results to the exclusion of those that don’t.

### 31.2.6. Insufficient Stack Depth

If the `errors` test results in a server crash at the `select infinite_recurse()` command, it means that the platform’s limit on process stack size is smaller than the `max_stack_depth` parameter indicates. This can be fixed by running the server under a higher stack size limit (4MB is recommended with the default value of `max_stack_depth`). If you are unable to do that, an alternative is to reduce the value of `max_stack_depth`.

On platforms supporting `getrlimit()`, the server should automatically choose a safe value of `max_stack_depth`; so unless you’ve manually overridden this setting, a failure of this kind is a reportable bug.

### 31.2.7. The “random” Test

The `random` test script is intended to produce random results. In very rare cases, this causes that regression test to fail. Typing:

```
diff results/random.out expected/random.out
```

should produce only one or a few lines of differences. You need not worry unless the random test fails repeatedly.

### 31.2.8. Configuration Parameters

When running the tests against an existing installation, some non-default parameter settings could cause the tests to fail. For example, changing parameters such as `enable_seqscan` or `enable_indexscan` could cause plan changes that would affect the results of tests that use `EXPLAIN`.

### 31.3. Variant Comparison Files

Since some of the tests inherently produce environment-dependent results, we have provided ways to specify alternate “expected” result files. Each regression test can have several comparison files showing possible results on different platforms. There are two independent mechanisms for determining which comparison file is used for each test.
The first mechanism allows comparison files to be selected for specific platforms. There is a mapping file, `src/test/regress/resultmap`, that defines which comparison file to use for each platform. To eliminate bogus test “failures” for a particular platform, you first choose or make a variant result file, and then add a line to the `resultmap` file.

Each line in the mapping file is of the form

```
testname:output:platformpattern=comparisonfilename
```

The test name is just the name of the particular regression test module. The output value indicates which output file to check. For the standard regression tests, this is always `out`. The value corresponds to the file extension of the output file. The platform pattern is a pattern in the style of the Unix tool `expr` (that is, a regular expression with an implicit `^` anchor at the start). It is matched against the platform name as printed by `config.guess`. The comparison file name is the base name of the substitute result comparison file.

For example: some systems interpret very small floating-point values as zero, rather than reporting an underflow error. This causes a few differences in the `float8` regression test. Therefore, we provide a variant comparison file, `float8-small-is-zero.out`, which includes the results to be expected on these systems. To silence the bogus “failure” message on OpenBSD platforms, `resultmap` includes:

```
float8:out:i.86-.*-openbsd=float8-small-is-zero.out
```

which will trigger on any machine where the output of `config.guess` matches `i.86-.*-openbsd.` Other lines in `resultmap` select the variant comparison file for other platforms where it’s appropriate.

The second selection mechanism for variant comparison files is much more automatic: it simply uses the “best match” among several supplied comparison files. The regression test driver script considers both the standard comparison file for a test, `testname.out`, and variant files named `testname_digit.out` (where the `digit` is any single digit `0-9`). If any such file is an exact match, the test is considered to pass; otherwise, the one that generates the shortest diff is used to create the failure report. (If `resultmap` includes an entry for the particular test, then the base `testname` is the substitute name given in `resultmap`.)

For example, for the `char` test, the comparison file `char.out` contains results that are expected in the `C` and `POSIX` locales, while the file `char_1.out` contains results sorted as they appear in many other locales.

The best-match mechanism was devised to cope with locale-dependent results, but it can be used in any situation where the test results cannot be predicted easily from the platform name alone. A limitation of this mechanism is that the test driver cannot tell which variant is actually “correct” for the current environment; it will just pick the variant that seems to work best. Therefore it is safest to use this mechanism only for variant results that you are willing to consider equally valid in all contexts.

### 31.4. TAP Tests

The client program tests under `src/bin` use the Perl TAP tools and are run by `prove`. You can pass command-line options to `prove` by setting the `make` variable `PROVE_FLAGS`, for example:

```
make -C src/bin check PROVE_FLAGS='--reverse'
```

The default is `--verbose`. See the manual page of `prove` for more information.
The tests written in Perl require the Perl module IPC::Run. This module is available from CPAN or an operating system package.

### 31.5. Test Coverage Examination

The PostgreSQL source code can be compiled with coverage testing instrumentation, so that it becomes possible to examine which parts of the code are covered by the regression tests or any other test suite that is run with the code. This is currently supported when compiling with GCC and requires the `gcov` and `lcov` programs.

A typical workflow would look like this:

```
./configure --enable-coverage ... OTHER OPTIONS ...
make
make check # or other test suite
make coverage-html
```

Then point your HTML browser to `coverage/index.html`. The `make` commands also work in subdirectories.

To reset the execution counts between test runs, run:

```
make coverage-clean
```
IV. Client Interfaces

This part describes the client programming interfaces distributed with PostgreSQL. Each of these chapters can be read independently. Note that there are many other programming interfaces for client programs that are distributed separately and contain their own documentation (Appendix H lists some of the more popular ones). Readers of this part should be familiar with using SQL commands to manipulate and query the database (see Part II) and of course with the programming language that the interface uses.
Chapter 32. libpq - C Library

libpq is the C application programmer’s interface to PostgreSQL. libpq is a set of library functions that allow client programs to pass queries to the PostgreSQL backend server and to receive the results of these queries.

libpq is also the underlying engine for several other PostgreSQL application interfaces, including those written for C++, Perl, Python, Tcl and ECPG. So some aspects of libpq’s behavior will be important to you if you use one of those packages. In particular, Section 32.14, Section 32.15 and Section 32.18 describe behavior that is visible to the user of any application that uses libpq.

Some short programs are included at the end of this chapter (Section 32.21) to show how to write programs that use libpq. There are also several complete examples of libpq applications in the directory src/test/examples in the source code distribution.

Client programs that use libpq must include the header file libpq-fe.h and must link with the libpq library.

32.1. Database Connection Control Functions

The following functions deal with making a connection to a PostgreSQL backend server. An application program can have several backend connections open at one time. (One reason to do that is to access more than one database.) Each connection is represented by a PGconn object, which is obtained from the function PQconnectdb, PQconnectdbParams, or PQsetdbLogin. Note that these functions will always return a non-null object pointer, unless perhaps there is too little memory even to allocate the PGconn object. The PQstatus function should be called to check the return value for a successful connection before queries are sent via the connection object.

Warning

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin each session by removing publicly-writable schemas from search_path. One can set parameter key word options to value -csearch_path=. Alternately, one can issue PQexec(conn, "SELECT pg_catalog.set_config('search_path', ", false)") after connecting. This consideration is not specific to libpq; it applies to every interface for executing arbitrary SQL commands.

Warning

On Unix, forking a process with open libpq connections can lead to unpredictable results because the parent and child processes share the same sockets and operating system resources. For this reason, such usage is not recommended, though doing an exec from the child process to load a new executable is safe.

Note: On Windows, there is a way to improve performance if a single database connection is repeatedly started and shutdown. Internally, libpq calls WSStartup() and WSACleanup() for connection startup and shutdown, respectively. WSStartup() increments an internal Windows library reference count which is decremented by WSACleanup(). When the reference count is just one, calling WSACleanup() frees all resources and all DLLs are unloaded. This is an expensive
operation. To avoid this, an application can manually call `WSAStartup()` so resources will not be freed when the last database connection is closed.

PQconnectdbParams

Makes a new connection to the database server.

```
PGconn *PQconnectdbParams(const char * const *keywords,
                            const char * const *values,
                            int expand_dbname);
```

This function opens a new database connection using the parameters taken from two NULL-terminated arrays. The first, `keywords`, is defined as an array of strings, each one being a key word. The second, `values`, gives the value for each key word. Unlike `PQsetdbLogin` below, the parameter set can be extended without changing the function signature, so use of this function (or its nonblocking analogs `PQconnectStartParams` and `PQconnectPoll`) is preferred for new application programming.

The currently recognized parameter key words are listed in Section 32.1.2.

When `expand_dbname` is non-zero, the `dbname` key word value is allowed to be recognized as a connection string. Only the first occurrence of `dbname` is expanded this way, any subsequent `dbname` value is processed as plain database name. More details on the possible connection string formats appear in Section 32.1.1.

The passed arrays can be empty to use all default parameters, or can contain one or more parameter settings. They should be matched in length. Processing will stop at the first NULL element in the `keywords` array.

If any parameter is NULL or an empty string, the corresponding environment variable (see Section 32.14) is checked. If the environment variable is not set either, then the indicated built-in defaults are used.

In general key words are processed from the beginning of these arrays in index order. The effect of this is that when key words are repeated, the last processed value is retained. Therefore, through careful placement of the `dbname` key word, it is possible to determine what may be overridden by a `conninfo` string, and what may not.

PQconnectdb

Makes a new connection to the database server.

```
PGconn *PQconnectdb(const char *conninfo);
```

This function opens a new database connection using the parameters taken from the string `conninfo`.

The passed string can be empty to use all default parameters, or it can contain one or more parameter settings separated by whitespace, or it can contain a URI. See Section 32.1.1 for details.

PQsetdbLogin

Makes a new connection to the database server.

```
PGconn *PQsetdbLogin(const char *pghost,
                     const char *pgport,
                     const char *pgoptions,
                     const char *pgtty,
                     const char *dbName,
                     const char *login,
                     const char *errorMessage);```
const char *pwd);

This is the predecessor of PQconnectdb with a fixed set of parameters. It has the same functionality except that the missing parameters will always take on default values. Write NULL or an empty string for any one of the fixed parameters that is to be defaulted.

If the dbName contains an = sign or has a valid connection URI prefix, it is taken as a conninfo string in exactly the same way as if it had been passed to PQconnectdb, and the remaining parameters are then applied as specified for PQconnectdbParams.

PQsetdb

Makes a new connection to the database server.

PGconn *PQsetdb(char *pghost,
    char *pgport,
    char *pgoptions,
    char *pgtty,
    char *dbName);

This is a macro that calls PQsetdbLogin with null pointers for the login and pwd parameters. It is provided for backward compatibility with very old programs.

PQconnectStartParams
PQconnectStart
PQconnectPoll

Make a connection to the database server in a nonblocking manner.

PGconn *PQconnectStartParams(const char * const *keywords,
    const char * const *values,
    int expand_dbname);

PGconn *PQconnectStart(const char *conninfo);

PostgresPollingStatusType PQconnectPoll(PGconn *conn);

These three functions are used to open a connection to a database server such that your application’s thread of execution is not blocked on remote I/O whilst doing so. The point of this approach is that the waits for I/O to complete can occur in the application’s main loop, rather than down inside PQconnectdbParams or PQconnectdb, and so the application can manage this operation in parallel with other activities.

With PQconnectStartParams, the database connection is made using the parameters taken from the keywords and values arrays, and controlled by expand_dbname, as described above for PQconnectdbParams.

With PQconnectStart, the database connection is made using the parameters taken from the string conninfo as described above for PQconnectdb.

Neither PQconnectStartParams nor PQconnectStart nor PQconnectPoll will block, so long as a number of restrictions are met:

- The hostaddr and host parameters are used appropriately to ensure that name and reverse name queries are not made. See the documentation of these parameters in Section 32.1.2 for details.
- If you call PQtrace, ensure that the stream object into which you trace will not block.
- You ensure that the socket is in the appropriate state before calling PQconnectPoll, as described below.
Note: use of PQconnectStartParams is analogous to PQconnectStart shown below.

To begin a nonblocking connection request, call

```c
conn = PQconnectStart("connection_info_string");
```

If `conn` is null, then libpq has been unable to allocate a new PGconn structure. Otherwise, a valid PGconn pointer is returned (though not yet representing a valid connection to the database). On return from PQconnectStart, call

```c
status = PQstatus(conn);
```

If `status` equals CONNECTION_BAD, PQconnectStart has failed.

If PQconnectStart succeeds, the next stage is to poll libpq so that it can proceed with the connection sequence. Use PQsocket(conn) to obtain the descriptor of the socket underlying the database connection. Loop thus:

- If PQconnectPoll(conn) last returned PGRES_POLLING_READING, wait until the socket is ready to read (as indicated by select(), poll(), or similar system function). Then call PQconnectPoll(conn) again.
- Conversely, if PQconnectPoll(conn) last returned PGRES_POLLING_WRITING, wait until the socket is ready to write, then call PQconnectPoll(conn) again. If you have yet to call PQconnectPoll, i.e., just after the call to PQconnectStart, behave as if it last returned PGRES_POLLING_WRITING. Continue this loop until PQconnectPoll(conn) returns PGRES_POLLING_FAILED, indicating the connection procedure has failed, or PGRES_POLLING_OK, indicating the connection has been successfully made.

At any time during connection, the status of the connection can be checked by calling PQstatus. If this call returns CONNECTION_BAD, then the connection procedure has failed; if the call returns CONNECTION_OK, then the connection is ready. Both of these states are equally detectable from the return value of PQconnectPoll, described above. Other states might also occur during (and only during) an asynchronous connection procedure. These indicate the current stage of the connection procedure and might be useful to provide feedback to the user for example. These statuses are:

- CONNECTION_STARTED
  Waiting for connection to be made.
- CONNECTION_MADE
  Connection OK; waiting to send.
- CONNECTION_AWAITING_RESPONSE
  Waiting for a response from the server.
- CONNECTION_AUTH_OK
  Received authentication; waiting for backend start-up to finish.
- CONNECTION_SSL_STARTUP
  Negotiating SSL encryption.
- CONNECTION_SETENV
  Negotiating environment-driven parameter settings.

Note that, although these constants will remain (in order to maintain compatibility), an application should never rely upon these occurring in a particular order, or at all, or on the status always being one of these documented values. An application might do something like this:

```c
switch(PQstatus(conn))
{
    case CONNECTION_STARTED:
        feedback = "Connecting...";
        break;
```
case CONNECTION_MADE:
    feedback = "Connected to server...";
    break;

.
.
.

default:
    feedback = "Connecting...";
}

The connect_timeout connection parameter is ignored when using PQconnectPoll; it is the application’s responsibility to decide whether an excessive amount of time has elapsed. Otherwise, PQconnectStart followed by a PQconnectPoll loop is equivalent to PQconnectdb.

Note that if PQconnectStart returns a non-null pointer, you must call PQfinish when you are finished with it, in order to dispose of the structure and any associated memory blocks. This must be done even if the connection attempt fails or is abandoned.

PQconndefaults

Returns the default connection options.

PQconninfoOption *PQconndefaults(void);

typedef struct
{
    char *keyword; /* The keyword of the option */
    char *envvar; /* Fallback environment variable name */
    char *compiled; /* Fallback compiled in default value */
    char *val; /* Option’s current value, or NULL */
    char *label; /* Label for field in connect dialog */
    char *dispchar; /* Indicates how to display this field in a connect dialog. Values are:
        ""     Display entered value as is
        "*"   Password field - hide value
        "D"   Debug option - don’t show by default */
    int dispsize; /* Field size in characters for dialog */
} PQconninfoOption;

Returns a connection options array. This can be used to determine all possible PQconnectdb options and their current default values. The return value points to an array of PQconninfoOption structures, which ends with an entry having a null keyword pointer. The null pointer is returned if memory could not be allocated. Note that the current default values (val fields) will depend on environment variables and other context. A missing or invalid service file will be silently ignored. Callers must treat the connection options data as read-only.

After processing the options array, free it by passing it to PQconninfoFree. If this is not done, a small amount of memory is leaked for each call to PQconndefaults.

PQconninfo

Returns the connection options used by a live connection.

PQconninfoOption *PQconninfo(PGconn *conn);

Returns a connection options array. This can be used to determine all possible PQconnectdb options and the values that were used to connect to the server. The return value points to an array of PQconninfoOption structures, which ends with an entry having a null keyword pointer. All notes above for PQconndefaults also apply to the result of PQconninfo.
PQconninfoParse

Returns parsed connection options from the provided connection string.

```c
PQconninfoOption *PQconninfoParse(const char *conninfo, char **errmsg);
```

 Parses a connection string and returns the resulting options as an array; or returns NULL if there is a problem with the connection string. This function can be used to extract the PQconnectdb options in the provided connection string. The return value points to an array of PQconninfoOption structures, which ends with an entry having a null keyword pointer.

All legal options will be present in the result array, but the PQconninfoOption for any option not present in the connection string will have val set to NULL; default values are not inserted.

If errmsg is not NULL, then *errmsg is set to NULL on success, else to a malloc’d error string explaining the problem. (It is also possible for *errmsg to be set to NULL and the function to return NULL; this indicates an out-of-memory condition.)

After processing the options array, free it by passing it to PQconninfoFree. If this is not done, some memory is leaked for each call to PQconninfoParse. Conversely, if an error occurs and errmsg is not NULL, be sure to free the error string using PQfreemem.

PQfinish

Closes the connection to the server. Also frees memory used by the PGconn object.

```c
void PQfinish(PGconn *conn);
```

Note that even if the server connection attempt fails (as indicated by PQstatus), the application should call PQfinish to free the memory used by the PGconn object. The PGconn pointer must not be used again after PQfinish has been called.

PQreset

Resets the communication channel to the server.

```c
void PQreset(PGconn *conn);
```

This function will close the connection to the server and attempt to reestablish a new connection to the same server, using all the same parameters previously used. This might be useful for error recovery if a working connection is lost.

PQresetStart

PQresetPoll

Reset the communication channel to the server, in a nonblocking manner.

```c
int PQresetStart(PGconn *conn);
```

```c
PostgresPollingStatusType PQresetPoll(PGconn *conn);
```

These functions will close the connection to the server and attempt to reestablish a new connection to the same server, using all the same parameters previously used. This can be useful for error recovery if a working connection is lost. They differ from PQreset (above) in that they act in a nonblocking manner. These functions suffer from the same restrictions as PQconnectStartParams, PQconnectStart and PQconnectPoll.

To initiate a connection reset, call PQresetStart. If it returns 0, the reset has failed. If it returns 1, poll the reset using PQresetPoll in exactly the same way as you would create the connection using PQconnectPoll.

PQpingParams

PQpingParams reports the status of the server. It accepts connection parameters identical to those of PQconnectdbParams, described above. It is not necessary to supply correct user name,
password, or database name values to obtain the server status; however, if incorrect values are
provided, the server will log a failed connection attempt.

```
PQPing PQpingParams(const char * const *keywords,
            const char * const *values,
            int expand_dbname);
```

The function returns one of the following values:

- **PQPING_OK**
  
  The server is running and appears to be accepting connections.

- **PQPING_REJECT**
  
  The server is running but is in a state that disallows connections (startup, shutdown, or crash
  recovery).

- **PQPING_NO_RESPONSE**
  
  The server could not be contacted. This might indicate that the server is not running, or that
  there is something wrong with the given connection parameters (for example, wrong port
  number), or that there is a network connectivity problem (for example, a firewall blocking
  the connection request).

- **PQPING_NO_ATTEMPT**
  
  No attempt was made to contact the server, because the supplied parameters were obviously
  incorrect or there was some client-side problem (for example, out of memory).

```
PQping
```

- **PQping** reports the status of the server. It accepts connection parameters identical to those of
  **PQconnectdb**, described above. It is not necessary to supply correct user name, password, or
database name values to obtain the server status; however, if incorrect values are provided, the
server will log a failed connection attempt.

```
PQPing PQping(const char *conninfo);
```

The return values are the same as for **PQPingParams**.

### 32.1.1. Connection Strings

Several libpq functions parse a user-specified string to obtain connection parameters. There are two
accepted formats for these strings: plain **keyword = value** strings and RFC 3986¹ URIs.

#### 32.1.1.1. Keyword/Value Connection Strings

In the first format, each parameter setting is in the form **keyword = value**. Spaces around the equal
sign are optional. To write an empty value, or a value containing spaces, surround it with single quotes,
e.g., **keyword = 'a value'**. Single quotes and backslashes within the value must be escaped with a
backslash, i.e., `'` and `\`.

Example:

```
host=localhost port=5432 dbname=mydb connect_timeout=10
```

---

The recognized parameter key words are listed in Section 32.1.2.

### 32.1.1.2. Connection URIs

The general form for a connection URI is:

```
postgresql://[user[:password]@][netloc][:port][/dbname][?param1=value1&...]
```

The URI scheme designator can be either `postgresql://` or `postgres://`. Each of the URI parts is optional. The following examples illustrate valid URI syntax uses:

- `postgresql://`
- `postgresql://localhost`
- `postgresql://localhost:5433`
- `postgresql://localhost/mydb`
- `postgresql://user@localhost`
- `postgresql://user:secret@localhost`
- `postgresql://other@localhost/otherdb?connect_timeout=10&application_name=myapp`

Components of the hierarchical part of the URI can also be given as parameters. For example:

```
postgresql:///mydb?host=localhost&port=5433
```

Percent-encoding may be used to include symbols with special meaning in any of the URI parts.

Any connection parameters not corresponding to key words listed in Section 32.1.2 are ignored and a warning message about them is sent to `stderr`.

For improved compatibility with JDBC connection URIs, instances of parameter `ssl=true` are translated into `sslmode=require`.

The host part may be either host name or an IP address. To specify an IPv6 host address, enclose it in square brackets:

```
postgresql://[2001:db8::1234]/database
```

The host component is interpreted as described for the parameter `host`. In particular, a Unix-domain socket connection is chosen if the host part is either empty or starts with a slash, otherwise a TCP/IP connection is initiated. Note, however, that the slash is a reserved character in the hierarchical part of the URI. So, to specify a non-standard Unix-domain socket directory, either omit the host specification in the URI and specify the host as a parameter, or percent-encode the path in the host component of the URI:

```
postgresql:///dbname?host=/var/lib/postgresql
postgresql://%2Fvar%2Flib%2Fpostgresql/dbname
```
32.1.2. Parameter Key Words

The currently recognized parameter key words are:

**host**
Name of host to connect to. If this begins with a slash, it specifies Unix-domain communication rather than TCP/IP communication; the value is the name of the directory in which the socket file is stored. The default behavior when host is not specified is to connect to a Unix-domain socket in /tmp (or whatever socket directory was specified when PostgreSQL was built). On machines without Unix-domain sockets, the default is to connect to localhost.

**hostaddr**
Numeric IP address of host to connect to. This should be in the standard IPv4 address format, e.g., 172.28.40.9. If your machine supports IPv6, you can also use those addresses. TCP/IP communication is always used when a nonempty string is specified for this parameter. Using hostaddr instead of host allows the application to avoid a host name look-up, which might be important in applications with time constraints. However, a host name is required for GSSAPI or SSPI authentication methods, as well as for verify-full SSL certificate verification. The following rules are used:

- If host is specified without hostaddr, a host name lookup occurs.
- If hostaddr is specified without host, the value for hostaddr gives the server network address. The connection attempt will fail if the authentication method requires a host name.
- If both host and hostaddr are specified, the value for hostaddr gives the server network address. The value for host is ignored unless the authentication method requires it, in which case it will be used as the host name.

Note that authentication is likely to fail if host is not the name of the server at network address hostaddr. Also, note that host rather than hostaddr is used to identify the connection in ~/.pgpass (see Section 32.15).

Without either a host name or host address, libpq will connect using a local Unix-domain socket; or on machines without Unix-domain sockets, it will attempt to connect to localhost.

**port**
Port number to connect to at the server host, or socket file name extension for Unix-domain connections.

**dbname**
The database name. Defaults to be the same as the user name. In certain contexts, the value is checked for extended formats; see Section 32.1.1 for more details on those.

**user**
PostgreSQL user name to connect as. Defaults to be the same as the operating system name of the user running the application.

**password**
Password to be used if the server demands password authentication.

**connect_timeout**
Maximum wait for connection, in seconds (write as a decimal integer string). Zero or not specified means wait indefinitely. It is not recommended to use a timeout of less than 2 seconds.
client_encoding

This sets the client_encoding configuration parameter for this connection. In addition to the values accepted by the corresponding server option, you can use auto to determine the right encoding from the current locale in the client (LC_CTYPE environment variable on Unix systems).

options

Specifies command-line options to send to the server at connection start. For example, setting this to -c geqo=off sets the session’s value of the geqo parameter to off. Spaces within this string are considered to separate command-line arguments, unless escaped with a backslash (\); write ```\``` to represent a literal backslash. For a detailed discussion of the available options, consult Chapter 19.

application_name

Specifies a value for the application_name configuration parameter.

fallback_application_name

Specifies a fallback value for the application_name configuration parameter. This value will be used if no value has been given for application_name via a connection parameter or the PGAPPNAME environment variable. Specifying a fallback name is useful in generic utility programs that wish to set a default application name but allow it to be overridden by the user.

keepalives

Controls whether client-side TCP keepalives are used. The default value is 1, meaning on, but you can change this to 0, meaning off, if keepalives are not wanted. This parameter is ignored for connections made via a Unix-domain socket.

keepalives_idle

Controls the number of seconds of inactivity after which TCP should send a keepalive message to the server. A value of zero uses the system default. This parameter is ignored for connections made via a Unix-domain socket, or if keepalives are disabled. It is only supported on systems where TCP_KEEPIDLE or an equivalent socket option is available, and on Windows; on other systems, it has no effect.

keepalives_interval

Controls the number of seconds after which a TCP keepalive message that is not acknowledged by the server should be retransmitted. A value of zero uses the system default. This parameter is ignored for connections made via a Unix-domain socket, or if keepalives are disabled. It is only supported on systems where TCP_KEEPINTVL or an equivalent socket option is available, and on Windows; on other systems, it has no effect.

keepalives_count

Controls the number of TCP keepalives that can be lost before the client’s connection to the server is considered dead. A value of zero uses the system default. This parameter is ignored for connections made via a Unix-domain socket, or if keepalives are disabled. It is only supported on systems where TCP_KEEPCNT or an equivalent socket option is available; on other systems, it has no effect.

tty

Ignored (formerly, this specified where to send server debug output).
sslmode

This option determines whether or with what priority a secure SSL TCP/IP connection will be negotiated with the server. There are six modes:

disable
  only try a non-SSL connection
allow
  first try a non-SSL connection; if that fails, try an SSL connection
prefer (default)
  first try an SSL connection; if that fails, try a non-SSL connection
require
  only try an SSL connection. If a root CA file is present, verify the certificate in the same way as if verify-ca was specified
verify-ca
  only try an SSL connection, and verify that the server certificate is issued by a trusted certificate authority (CA)
verify-full
  only try an SSL connection, verify that the server certificate is issued by a trusted CA and that the requested server host name matches that in the certificate

See Section 32.18 for a detailed description of how these options work.

sslmode is ignored for Unix domain socket communication. If PostgreSQL is compiled without SSL support, using options require, verify-ca, or verify-full will cause an error, while options allow and prefer will be accepted but libpq will not actually attempt an SSL connection.

requiressl

This option is deprecated in favor of the sslmode setting.

If set to 1, an SSL connection to the server is required (this is equivalent to sslmode require). libpq will then refuse to connect if the server does not accept an SSL connection. If set to 0 (default), libpq will negotiate the connection type with the server (equivalent to sslmode prefer). This option is only available if PostgreSQL is compiled with SSL support.

sslcompression

If set to 1 (default), data sent over SSL connections will be compressed (this requires OpenSSL version 0.9.8 or later). If set to 0, compression will be disabled (this requires OpenSSL 1.0.0 or later). This parameter is ignored if a connection without SSL is made, or if the version of OpenSSL used does not support it.

Compression uses CPU time, but can improve throughput if the network is the bottleneck. Disabling compression can improve response time and throughput if CPU performance is the limiting factor.

sslcert

This parameter specifies the file name of the client SSL certificate, replacing the default ~/.postgresql/postgresql.crt. This parameter is ignored if an SSL connection is not made.
sslkey

This parameter specifies the location for the secret key used for the client certificate. It can either specify a file name that will be used instead of the default ~/.postgresql/postgresql.key, or it can specify a key obtained from an external “engine” (engines are OpenSSL loadable modules). An external engine specification should consist of a colon-separated engine name and an engine-specific key identifier. This parameter is ignored if an SSL connection is not made.

sslrootcert

This parameter specifies the name of a file containing SSL certificate authority (CA) certificate(s). If the file exists, the server’s certificate will be verified to be signed by one of these authorities. The default is ~/.postgresql/root.crt.

sslcrl

This parameter specifies the file name of the SSL certificate revocation list (CRL). Certificates listed in this file, if it exists, will be rejected while attempting to authenticate the server’s certificate. The default is ~/.postgresql/root.crl.

requirepeer

This parameter specifies the operating-system user name of the server, for example requirepeer=postgres. When making a Unix-domain socket connection, if this parameter is set, the client checks at the beginning of the connection that the server process is running under the specified user name; if it is not, the connection is aborted with an error. This parameter can be used to provide server authentication similar to that available with SSL certificates on TCP/IP connections. (Note that if the Unix-domain socket is in /tmp or another publicly writable location, any user could start a server listening there. Use this parameter to ensure that you are connected to a server run by a trusted user.) This option is only supported on platforms for which the peer authentication method is implemented; see Section 20.3.6.

krbsrvname

Kerberos service name to use when authenticating with GSSAPI. This must match the service name specified in the server configuration for Kerberos authentication to succeed. (See also Section 20.3.3.)

gsslib

GSS library to use for GSSAPI authentication. Currently this is disregarded except on Windows builds that include both GSSAPI and SSPI support. In that case, set this to gssapi to cause libpq to use the GSSAPI library for authentication instead of the default SSPI.

service

Service name to use for additional parameters. It specifies a service name in pg_service.conf that holds additional connection parameters. This allows applications to specify only a service name so connection parameters can be centrally maintained. See Section 32.16.

32.2. Connection Status Functions

These functions can be used to interrogate the status of an existing database connection object.

Tip: libpq application programmers should be careful to maintain the PGconn abstraction. Use the accessor functions described below to get at the contents of PGconn. Reference to internal
The following functions return parameter values established at connection. These values are fixed for the life of the PGconn object.

PQdb

Returns the database name of the connection.
char *PQdb(const PGconn *conn);

PQuser

Returns the user name of the connection.
char *PQuser(const PGconn *conn);

PQpass

Returns the password of the connection.
char *PQpass(const PGconn *conn);

PQhost

Returns the server host name of the connection. This can be a host name, an IP address, or a directory path if the connection is via Unix socket. (The path case can be distinguished because it will always be an absolute path, beginning with /.)
char *PQhost(const PGconn *conn);

PQport

Returns the port of the connection.
char *PQport(const PGconn *conn);

PQtty

Returns the debug TTY of the connection. (This is obsolete, since the server no longer pays attention to the TTY setting, but the function remains for backward compatibility.)
char *PQtty(const PGconn *conn);

PQoptions

Returns the command-line options passed in the connection request.
char *PQoptions(const PGconn *conn);

The following functions return status data that can change as operations are executed on the PGconn object.

PQstatus

Returns the status of the connection.
ConnStatusType PQstatus(const PGconn *conn);

The status can be one of a number of values. However, only two of these are seen outside of an asynchronous connection procedure: CONNECTION_OK and CONNECTION_BAD. A good connection to the database has the status CONNECTION_OK. A failed connection attempt is signaled by
status CONNECTION_BAD. Ordinarily, an OK status will remain so until PQfinish, but a communications failure might result in the status changing to CONNECTION_BAD prematurely. In that case the application could try to recover by calling PQreset.

See the entry for PQconnectStartParams, PQconnectStart and PQconnectPoll with regards to other status codes that might be returned.

PQtransactionStatus

Returns the current in-transaction status of the server.

PGTransactionStatusType PQtransactionStatus(const PGconn *conn);
The status can be PQTRANS_IDLE (currently idle), PQTRANS_ACTIVE (a command is in progress), PQTRANS_INTRTRAN (idle, in a valid transaction block), or PQTRANS_INERROR (idle, in a failed transaction block). PQTRANS_UNKNOWN is reported if the connection is bad. PQTRANS_ACTIVE is reported only when a query has been sent to the server and not yet completed.

PQparameterStatus

Looks up a current parameter setting of the server.

const char *PQparameterStatus(const PGconn *conn, const char *paramName);
Certain parameter values are reported by the server automatically at connection startup or whenever their values change. PQparameterStatus can be used to interrogate these settings. It returns the current value of a parameter if known, or NULL if the parameter is not known.

Parameters reported as of the current release include server_version, server_encoding, client_encoding, application_name, is_superuser, session_authorization, DateStyle, IntervalStyle, TimeZone, integer_datetimes, and standard_conforming_strings. (server_encoding, TimeZone, and integer_datetimes were not reported by releases before 8.0; standard_conforming_strings was not reported by releases before 8.1; IntervalStyle was not reported by releases before 8.4; application_name was not reported by releases before 9.0.) Note that server_version, server_encoding and integer_datetimes cannot change after startup.

Pre-3.0-protocol servers do not report parameter settings, but libpq includes logic to obtain values for server_version and client_encoding anyway. Applications are encouraged to use PQparameterStatus rather than ad hoc code to determine these values. (Beware however that on a pre-3.0 connection, changing client_encoding via SET after connection startup will not be reflected by PQparameterStatus.) For server_version, see also PQserverVersion, which returns the information in a numeric form that is much easier to compare against.

If no value for standard_conforming_strings is reported, applications can assume it is off, that is, backslashes are treated as escapes in string literals. Also, the presence of this parameter can be taken as an indication that the escape string syntax (E’...’) is accepted.

Although the returned pointer is declared const, it in fact points to mutable storage associated with the PGconn structure. It is unwise to assume the pointer will remain valid across queries.

PQprotocolVersion

Interrogates the frontend/backend protocol being used.

int PQprotocolVersion(const PGconn *conn);
Applications might wish to use this function to determine whether certain features are supported. Currently, the possible values are 2 (2.0 protocol), 3 (3.0 protocol), or zero (connection bad). The protocol version will not change after connection startup is complete, but it could theoretically change during a connection reset. The 3.0 protocol will normally be used when communicating
with PostgreSQL 7.4 or later servers; pre-7.4 servers support only protocol 2.0. (Protocol 1.0 is obsolete and not supported by libpq.)

PQserverVersion

Returns an integer representing the backend version.

```c
int PQserverVersion(const PGconn *conn);
```

Applications might use this function to determine the version of the database server they are connected to. The number is formed by converting the major, minor, and revision numbers into two-decimal-digit numbers and appending them together. For example, version 8.1.5 will be returned as 80105, and version 8.2 will be returned as 80200 (leading zeroes are not shown). Zero is returned if the connection is bad.

PQerrorMessage

Returns the error message most recently generated by an operation on the connection.

```c
char *PQerrorMessage(const PGconn *conn);
```

Nearly all libpq functions will set a message for PQerrorMessage if they fail. Note that by libpq convention, a nonempty PQerrorMessage result can consist of multiple lines, and will include a trailing newline. The caller should not free the result directly. It will be freed when the associated PGconn handle is passed to PQfinish. The result string should not be expected to remain the same across operations on the PGconn structure.

PQsocket

Obtains the file descriptor number of the connection socket to the server. A valid descriptor will be greater than or equal to 0; a result of -1 indicates that no server connection is currently open. (This will not change during normal operation, but could change during connection setup or reset.)

```c
int PQsocket(const PGconn *conn);
```

PQbackendPID

Returns the process ID (PID) of the backend process handling this connection.

```c
int PQbackendPID(const PGconn *conn);
```

The backend PID is useful for debugging purposes and for comparison to NOTIFY messages (which include the PID of the notifying backend process). Note that the PID belongs to a process executing on the database server host, not the local host!

PQconnectionNeedsPassword

Returns true (1) if the connection authentication method required a password, but none was available. Returns false (0) if not.

```c
int PQconnectionNeedsPassword(const PGconn *conn);
```

This function can be applied after a failed connection attempt to decide whether to prompt the user for a password.

PQconnectionUsedPassword

Returns true (1) if the connection authentication method used a password. Returns false (0) if not.

```c
int PQconnectionUsedPassword(const PGconn *conn);
```

This function can be applied after either a failed or successful connection attempt to detect whether the server demanded a password.
The following functions return information related to SSL. This information usually doesn’t change after a connection is established.

PQsslInUse

Returns true (1) if the connection uses SSL, false (0) if not.

int PQsslInUse(const PGconn *conn);

PQsslAttribute

Returns SSL-related information about the connection.

const char *PQsslAttribute(const PGconn *conn, const char *attribute_name);

The list of available attributes varies depending on the SSL library being used, and the type of connection. If an attribute is not available, returns NULL.

The following attributes are commonly available:

library

Name of the SSL implementation in use. (Currently, only "OpenSSL" is implemented)

protocol

SSL/TLS version in use. Common values are "SSLv2", "SSLv3", "TLSv1", "TLSv1.1" and "TLSv1.2", but an implementation may return other strings if some other protocol is used.

key_bits

Number of key bits used by the encryption algorithm.

cipher

A short name of the ciphersuite used, e.g. "DHE-RSA-DES-CBC3-SHA". The names are specific to each SSL implementation.

compression

If SSL compression is in use, returns the name of the compression algorithm, or "on" if compression is used but the algorithm is not known. If compression is not in use, returns "off".

PQsslAttributeNames

Return an array of SSL attribute names available. The array is terminated by a NULL pointer.

const char * const * PQsslAttributeNames(const PGconn *conn);

PQsslStruct

Return a pointer to an SSL-implementation-specific object describing the connection.

void *PQsslStruct(const PGconn *conn, const char *struct_name);

The struct(s) available depend on the SSL implementation in use. For OpenSSL, there is one struct, available under the name "OpenSSL", and it returns a pointer to the OpenSSL SSL struct. To use this function, code along the following lines could be used:

```
#include <libpq-fe.h>
#include <openssl/ssl.h>

...```
SSL *ssl;

dbconn = PQconnectdb(...);
...

ssl = PQsslStruct(dbconn, "OpenSSL");
if (ssl)
{
    /* use OpenSSL functions to access ssl */
}

This structure can be used to verify encryption levels, check server certificates, and more. Refer to the OpenSSL documentation for information about this structure.

PQgetssl

Returns the SSL structure used in the connection, or null if SSL is not in use.

void *PQgetssl(const PGconn *conn);

This function is equivalent to PQsslStruct(conn, "OpenSSL"). It should not be used in new applications, because the returned struct is specific to OpenSSL and will not be available if another SSL implementation is used. To check if a connection uses SSL, call PQsslInUse instead, and for more details about the connection, use PQsslAttribute.

32.3. Command Execution Functions

Once a connection to a database server has been successfully established, the functions described here are used to perform SQL queries and commands.

32.3.1. Main Functions

PQexec

Submits a command to the server and waits for the result.

PGresult *PQexec(PGconn *conn, const char *command);

Returns a PGresult pointer or possibly a null pointer. A non-null pointer will generally be returned except in out-of-memory conditions or serious errors such as inability to send the command to the server. The PGresultStatus function should be called to check the return value for any errors (including the value of a null pointer, in which case it will return PGRES_FATAL_ERROR). Use PQerrorMessage to get more information about such errors.

The command string can include multiple SQL commands (separated by semicolons). Multiple queries sent in a single PQexec call are processed in a single transaction, unless there are explicit BEGIN/COMMIT commands included in the query string to divide it into multiple transactions. Note however that the returned PGresult structure describes only the result of the last command executed from the string. Should one of the commands fail, processing of the string stops with it and the returned PGresult describes the error condition.
PQexecParams

Submits a command to the server and waits for the result, with the ability to pass parameters separately from the SQL command text.

```
PQresult *PQexecParams(PGconn *conn, const char *command, int nParams, const Oid *paramTypes, const char * const *paramValues, const int *paramLengths, const int *paramFormats, int resultFormat);
```

PQexecParams is like PQexec, but offers additional functionality: parameter values can be specified separately from the command string proper, and query results can be requested in either text or binary format. PQexecParams is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

The function arguments are:

- `conn` The connection object to send the command through.
- `command` The SQL command string to be executed. If parameters are used, they are referred to in the command string as $1, $2, etc.
- `nParams` The number of parameters supplied; it is the length of the arrays `paramTypes[]`, `paramValues[]`, `paramLengths[]`, and `paramFormats[]`. (The array pointers can be NULL when `nParams` is zero.)
- `paramTypes[]` Specifies, by OID, the data types to be assigned to the parameter symbols. If `paramTypes` is NULL, or any particular element in the array is zero, the server infers a data type for the parameter symbol in the same way it would do for an untyped literal string.
- `paramValues[]` Specifies the actual values of the parameters. A null pointer in this array means the corresponding parameter is null; otherwise the pointer points to a zero-terminated text string (for text format) or binary data in the format expected by the server (for binary format).
- `paramLengths[]` Specifies the actual data lengths of binary-format parameters. It is ignored for null parameters and text-format parameters. The array pointer can be null when there are no binary parameters.
- `paramFormats[]` Specifies whether parameters are text (put a zero in the array entry for the corresponding parameter) or binary (put a one in the array entry for the corresponding parameter). If the array pointer is null then all parameters are presumed to be text strings.

Values passed in binary format require knowledge of the internal representation expected by the backend. For example, integers must be passed in network byte order. Passing numeric values requires knowledge of the server storage format, as
implemented in `src/backend/utils/adt/numeric.c::numeric_send()` and `src/backend/utils/adt/numeric.c::numeric_recv()`.

**resultFormat**

Specify zero to obtain results in text format, or one to obtain results in binary format. (There is not currently a provision to obtain different result columns in different formats, although that is possible in the underlying protocol.)

The primary advantage of `PQexecParams` over `PQexec` is that parameter values can be separated from the command string, thus avoiding the need for tedious and error-prone quoting and escaping. Unlike `PQexec`, `PQexecParams` allows at most one SQL command in the given string. (There can be semicolons in it, but not more than one nonempty command.) This is a limitation of the underlying protocol, but has some usefulness as an extra defense against SQL-injection attacks.

**Tip:** Specifying parameter types via OIDs is tedious, particularly if you prefer not to hard-wire particular OID values into your program. However, you can avoid doing so even in cases where the server by itself cannot determine the type of the parameter, or chooses a different type than you want. In the SQL command text, attach an explicit cast to the parameter symbol to show what data type you will send. For example:

```
SELECT * FROM mytable WHERE x = $1::bigint;
```

This forces parameter $1 to be treated as bigint, whereas by default it would be assigned the same type as x. Forcing the parameter type decision, either this way or by specifying a numeric type OID, is strongly recommended when sending parameter values in binary format, because binary format has less redundancy than text format and so there is less chance that the server will detect a type mismatch mistake for you.

**PQprepare**

Submits a request to create a prepared statement with the given parameters, and waits for completion.

```
PQresult *PQprepare(PGconn *conn, const char *stmtName, const char *query, int nParams, const Oid *paramTypes);
```

`PQprepare` creates a prepared statement for later execution with `PQexecPrepared`. This feature allows commands to be executed repeatedly without being parsed and planned each time; see `PREPARE` for details. `PQprepare` is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

The function creates a prepared statement named `stmtName` from the `query` string, which must contain a single SQL command. `stmtName` can be "" to create an unnamed statement, in which case any pre-existing unnamed statement is automatically replaced; otherwise it is an error if the statement name is already defined in the current session. If any parameters are used, they are referred to in the query as $1, $2, etc. `nParams` is the number of parameters for which types are pre-specified in the array `paramTypes[]`. (The array pointer can be NULL when `nParams` is zero.) `paramTypes[]` specifies, by OID, the data types to be assigned to the parameter symbols.
If `paramTypes` is `NULL`, or any particular element in the array is zero, the server assigns a data type to the parameter symbol in the same way it would do for an untyped literal string. Also, the query can use parameter symbols with numbers higher than `nParams`; data types will be inferred for these symbols as well. (See `PQdescribePrepared` for a means to find out what data types were inferred.)

As with `PQexec`, the result is normally a `PGresult` object whose contents indicate server-side success or failure. A null result indicates out-of-memory or inability to send the command at all. Use `PQerrorMessage` to get more information about such errors.

Prepared statements for use with `PQexecPrepared` can also be created by executing SQL PREPARE statements. Also, although there is no libpq function for deleting a prepared statement, the SQL DEALLOCATE statement can be used for that purpose.

`PQexecPrepared` sends a request to execute a prepared statement with given parameters, and waits for the result.

```c
PGresult *PQexecPrepared(PGconn *conn, const char *stmtName, int nParams, const char * const *paramValues, const int *paramLengths, const int *paramFormats, int resultFormat);
```

`PQexecPrepared` is like `PQexecParams`, but the command to be executed is specified by naming a previously-prepared statement, instead of giving a query string. This feature allows commands that will be used repeatedly to be parsed and planned just once, rather than each time they are executed. The statement must have been prepared previously in the current session. `PQexecPrepared` is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

The parameters are identical to `PQexecParams`, except that the name of a prepared statement is given instead of a query string, and the `paramTypes[]` parameter is not present (it is not needed since the prepared statement’s parameter types were determined when it was created).

`PQdescribePrepared` submits a request to obtain information about the specified prepared statement, and waits for completion.

```c
PGresult *PQdescribePrepared(PGconn *conn, const char *stmtName);
```

`PQdescribePrepared` allows an application to obtain information about a previously prepared statement. `PQdescribePrepared` is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

`stmtName` can be `""` or `NULL` to reference the unnamed statement, otherwise it must be the name of an existing prepared statement. On success, a `PGresult` with status `PGRES_COMMAND_OK` is returned. The functions `PQnparams` and `PQparamtype` can be applied to this `PGresult` to obtain information about the parameters of the prepared statement, and the functions `PQnfields`, `PQfname`, `PQftype`, etc provide information about the result columns (if any) of the statement.

`PQdescribePortal` submits a request to obtain information about the specified portal, and waits for completion.

```c
PGresult *PQdescribePortal(PGconn *conn, const char *portalName);
```

`PQdescribePortal` allows an application to obtain information about a previously created portal. (libpq does not provide any direct access to portals, but you can use this function to inspect the
properties of a cursor created with a DECLARE CURSOR SQL command.) PQdescribePortal is supported only in protocol 3.0 and later connections; it will fail when using protocol 2.0.

portalName can be "" or NULL to reference the unnamed portal, otherwise it must be the name of an existing portal. On success, a PGresult with status PGRES_COMMAND_OK is returned. The functions PQnfields, PQfname, PQftype, etc can be applied to the PGresult to obtain information about the result columns (if any) of the portal.

The PGresult structure encapsulates the result returned by the server. libpq application programmers should be careful to maintain the PGresult abstraction. Use the accessor functions below to get at the contents of PGresult. Avoid directly referencing the fields of the PGresult structure because they are subject to change in the future.

PQresultStatus

Returns the result status of the command.

ExecStatusType PQresultStatus(const PGresult *res);

PQresultStatus can return one of the following values:

PGRES_EMPTY_QUERY

The string sent to the server was empty.

PGRES_COMMAND_OK

Successful completion of a command returning no data.

PGRES_TUPLES_OK

Successful completion of a command returning data (such as a SELECT or SHOW).

PGRES_COPY_OUT

Copy Out (from server) data transfer started.

PGRES_COPY_IN

Copy In (to server) data transfer started.

PGRES_BAD_RESPONSE

The server’s response was not understood.

PGRES_NONFATAL_ERROR

A nonfatal error (a notice or warning) occurred.

PGRES_FATAL_ERROR

A fatal error occurred.

PGRES_COPY_BOTH

Copy In/Out (to and from server) data transfer started. This feature is currently used only for streaming replication, so this status should not occur in ordinary applications.

PGRES_SINGLE_TUPLE

The PGresult contains a single result tuple from the current command. This status occurs only when single-row mode has been selected for the query (see Section 32.5).

If the result status is PGRES_TUPLES_OK or PGRES_SINGLE_TUPLE, then the functions described below can be used to retrieve the rows returned by the query. Note that a SELECT command that happens to retrieve zero rows still shows PGRES_TUPLES_OK, PGRES_COMMAND_OK
is for commands that can never return rows (\texttt{INSERT} or \texttt{UPDATE} without a \texttt{RETURNING} clause, etc.). A response of \texttt{PGRES\_EMPTY\_QUERY} might indicate a bug in the client software.

A result of status \texttt{PGRES\_NONFATAL\_ERROR} will never be returned directly by \texttt{PQexec} or other query execution functions; results of this kind are instead passed to the notice processor (see Section 32.12).

\textbf{PQresStatus}

Converts the enumerated type returned by \texttt{PQresultStatus} into a string constant describing the status code. The caller should not free the result.

\begin{verbatim}
char *PQresStatus(ExecStatusType status);
\end{verbatim}

\textbf{PQresultErrorMessage}

Returns the error message associated with the command, or an empty string if there was no error.

\begin{verbatim}
char *PQresultErrorMessage(const PGresult *res);
\end{verbatim}

If there was an error, the returned string will include a trailing newline. The caller should not free the result directly. It will be freed when the associated \texttt{PGresult} handle is passed to \texttt{PQclear}.

Immediately following a \texttt{PQexec} or \texttt{PQgetResult} call, \texttt{PQerrorMessage} (on the connection) will return the same string as \texttt{PQresultErrorMessage} (on the result). However, a \texttt{PGresult} will retain its error message until destroyed, whereas the connection’s error message will change when subsequent operations are done. Use \texttt{PQresultErrorMessage} when you want to know the status associated with a particular \texttt{PGresult}; use \texttt{PQerrorMessage} when you want to know the status from the latest operation on the connection.

\textbf{PQresultVerboseErrorMessage}

Returns a reformatted version of the error message associated with a \texttt{PGresult} object.

\begin{verbatim}
char *PQresultVerboseErrorMessage(const PGresult *res, PGVerbosity verbosity, PGContextVisibility show_context);
\end{verbatim}

In some situations a client might wish to obtain a more detailed version of a previously-reported error. \texttt{PQresultVerboseErrorMessage} addresses this need by computing the message that would have been produced by \texttt{PQresultErrorMessage} if the specified verbosity settings had been in effect for the connection when the given \texttt{PGresult} was generated. If the \texttt{PGresult} is not an error result, “\texttt{PGresult is not an error result}” is reported instead. The returned string includes a trailing newline.

Unlike most other functions for extracting data from a \texttt{PGresult}, the result of this function is a freshly allocated string. The caller must free it using \texttt{PQfreemem()} when the string is no longer needed.

A NULL return is possible if there is insufficient memory.

\textbf{PQresultErrorField}

Returns an individual field of an error report.

\begin{verbatim}
char *PQresultErrorField(const PGresult *res, int fieldcode);
\end{verbatim}

\texttt{fieldcode} is an error field identifier; see the symbols listed below. NULL is returned if the \texttt{PGresult} is not an error or warning result, or does not include the specified field. Field values will normally not include a trailing newline. The caller should not free the result directly. It will be freed when the associated \texttt{PGresult} handle is passed to \texttt{PQclear}.

The following field codes are available:
The severity; the field contents are \texttt{ERROR}, \texttt{FATAL}, or \texttt{PANIC} (in an error message), or \texttt{WARNING}, \texttt{NOTICE}, \texttt{DEBUG}, \texttt{INFO}, or \texttt{LOG} (in a notice message), or a localized translation of one of these. Always present.

The severity; the field contents are \texttt{ERROR}, \texttt{FATAL}, or \texttt{PANIC} (in an error message), or \texttt{WARNING}, \texttt{NOTICE}, \texttt{DEBUG}, \texttt{INFO}, or \texttt{LOG} (in a notice message). This is identical to the \texttt{PG\_DIAG\_SEVERITY} field except that the contents are never localized. This is present only in reports generated by PostgreSQL versions 9.6 and later.

The SQLSTATE code for the error. The SQLSTATE code identifies the type of error that has occurred; it can be used by front-end applications to perform specific operations (such as error handling) in response to a particular database error. For a list of the possible SQLSTATE codes, see Appendix A. This field is not localizable, and is always present.

The primary human-readable error message (typically one line). Always present.

Detail: an optional secondary error message carrying more detail about the problem. Might run to multiple lines.

Hint: an optional suggestion what to do about the problem. This is intended to differ from detail in that it offers advice (potentially inappropriate) rather than hard facts. Might run to multiple lines.

A string containing a decimal integer indicating an error cursor position as an index into the original statement string. The first character has index 1, and positions are measured in characters not bytes.

This is defined the same as the \texttt{PG\_DIAG\_STATEMENT\_POSITION} field, but it is used when the cursor position refers to an internally generated command rather than the one submitted by the client. The \texttt{PG\_DIAG\_INTERNAL\_QUERY} field will always appear when this field appears.

The text of a failed internally-generated command. This could be, for example, a SQL query issued by a PL/pgSQL function.

An indication of the context in which the error occurred. Presently this includes a call stack traceback of active procedural language functions and internally-generated queries. The trace is one entry per line, most recent first.

If the error was associated with a specific database object, the name of the schema containing that object, if any.
Chapter 32. libpq - C Library

If the error was associated with a specific table, the name of the table. (Refer to the schema name field for the name of the table’s schema.)

If the error was associated with a specific table column, the name of the column. (Refer to the schema and table name fields to identify the table.)

If the error was associated with a specific data type, the name of the data type. (Refer to the schema name field for the name of the data type’s schema.)

If the error was associated with a specific constraint, the name of the constraint. Refer to fields listed above for the associated table or domain. (For this purpose, indexes are treated as constraints, even if they weren’t created with constraint syntax.)

The file name of the source-code location where the error was reported.

The line number of the source-code location where the error was reported.

The name of the source-code function reporting the error.

Note: The fields for schema name, table name, column name, data type name, and constraint name are supplied only for a limited number of error types; see Appendix A. Do not assume that the presence of any of these fields guarantees the presence of another field. Core error sources observe the interrelationships noted above, but user-defined functions may use these fields in other ways. In the same vein, do not assume that these fields denote contemporary objects in the current database.

The client is responsible for formatting displayed information to meet its needs; in particular it should break long lines as needed. Newline characters appearing in the error message fields should be treated as paragraph breaks, not line breaks.

Errors generated internally by libpq will have severity and primary message, but typically no other fields. Errors returned by a pre-3.0-protocol server will include severity and primary message, and sometimes a detail message, but no other fields.

Note that error fields are only available from PGresult objects, not PGconn objects; there is no PQerrorField function.

Frees the storage associated with a PGresult. Every command result should be freed via PQclear when it is no longer needed.

void PQclear(PGresult *res);

You can keep a PGresult object around for as long as you need it; it does not go away when you issue a new command, nor even if you close the connection. To get rid of it, you must call PQclear. Failure to do this will result in memory leaks in your application.
32.3.2. Retrieving Query Result Information

These functions are used to extract information from a `PGresult` object that represents a successful query result (that is, one that has status `PGRES_TUPLES_OK` or `PGRES_SINGLE_TUPLE`). They can also be used to extract information from a successful Describe operation: a Describe’s result has all the same column information that actual execution of the query would provide, but it has zero rows. For objects with other status values, these functions will act as though the result has zero rows and zero columns.

`PQntuples`

Returns the number of rows (tuples) in the query result. (Note that `PGresult` objects are limited to no more than `INT_MAX` rows, so an `int` result is sufficient.)

```c
int PQntuples(const PGresult *res);
```

`PQnfields`

Returns the number of columns (fields) in each row of the query result.

```c
int PQnfields(const PGresult *res);
```

`PQfname`

Returns the column name associated with the given column number. Column numbers start at 0. The caller should not free the result directly. It will be freed when the associated `PGresult` handle is passed to `PQclear`.

```c
char *PQfname(const PGresult *res,
               int column_number);
```

NULL is returned if the column number is out of range.

`PQfnumber`

Returns the column number associated with the given column name.

```c
int PQfnumber(const PGresult *res,
               const char *column_name);
```

-1 is returned if the given name does not match any column.

The given name is treated like an identifier in an SQL command, that is, it is downcased unless double-quoted. For example, given a query result generated from the SQL command:

```
SELECT 1 AS FOO, 2 AS "BAR";
```

we would have the results:

```
PQfname(res, 0)  foo
PQfname(res, 1)  BAR
PQfnumber(res, "FOO")  0
PQfnumber(res, "foo")  0
PQfnumber(res, "BAR")  -1
PQfnumber(res, "\"BAR\"")  1
```

`PQftable`

Returns the OID of the table from which the given column was fetched. Column numbers start at 0.

```c
Oid PQftable(const PGresult *res,
             int column_number);
```
InvalidOid is returned if the column number is out of range, or if the specified column is not a simple reference to a table column, or when using pre-3.0 protocol. You can query the system table pg_class to determine exactly which table is referenced.

The type Oid and the constant InvalidOid will be defined when you include the libpq header file. They will both be some integer type.

PQftablecol

Returns the column number (within its table) of the column making up the specified query result column. Query-result column numbers start at 0, but table columns have nonzero numbers.

```c
int PQftablecol(const PGresult *res, int column_number);
```

Zero is returned if the column number is out of range, or if the specified column is not a simple reference to a table column, or when using pre-3.0 protocol.

PQfformat

Returns the format code indicating the format of the given column. Column numbers start at 0.

```c
int PQfformat(const PGresult *res, int column_number);
```

Format code zero indicates textual data representation, while format code one indicates binary representation. (Other codes are reserved for future definition.)

PQftype

Returns the data type associated with the given column number. The integer returned is the internal OID number of the type. Column numbers start at 0.

```c
Oid PQftype(const PGresult *res, int column_number);
```

You can query the system table pg_type to obtain the names and properties of the various data types. The OIDs of the built-in data types are defined in the file src/include/catalog/pg_type.h in the source tree.

PQfmod

Returns the type modifier of the column associated with the given column number. Column numbers start at 0.

```c
int PQfmod(const PGresult *res, int column_number);
```

The interpretation of modifier values is type-specific; they typically indicate precision or size limits. The value -1 is used to indicate “no information available”. Most data types do not use modifiers, in which case the value is always -1.

PQfsize

Returns the size in bytes of the column associated with the given column number. Column numbers start at 0.

```c
int PQfsize(const PGresult *res, int column_number);
```

PQfsize returns the space allocated for this column in a database row, in other words the size of the server’s internal representation of the data type. (Accordingly, it is not really very useful to clients.) A negative value indicates the data type is variable-length.
PQbinaryTuples

Returns 1 if the PGresult contains binary data and 0 if it contains text data.

int PQbinaryTuples(const PGresult *res);

This function is deprecated (except for its use in connection with COPY), because it is possible for a single PGresult to contain text data in some columns and binary data in others. PQfformat is preferred. PQbinaryTuples returns 1 only if all columns of the result are binary (format 1).

PQgetvalue

Returns a single field value of one row of a PGresult. Row and column numbers start at 0. The caller should not free the result directly. It will be freed when the associated PGresult handle is passed to PQclear.

char *PQgetvalue(const PGresult *res,
                   int row_number,
                   int column_number);

For data in text format, the value returned by PQgetvalue is a null-terminated character string representation of the field value. For data in binary format, the value is in the binary representation determined by the data type’s typsend and typreceive functions. (The value is actually followed by a zero byte in this case too, but that is not ordinarily useful, since the value is likely to contain embedded nulls.)

An empty string is returned if the field value is null. See PQgetisnull to distinguish null values from empty-string values.

The pointer returned by PQgetvalue points to storage that is part of the PGresult structure. One should not modify the data it points to, and one must explicitly copy the data into other storage if it is to be used past the lifetime of the PGresult structure itself.

PQgetisnull

Tests a field for a null value. Row and column numbers start at 0.

int PQgetisnull(const PGresult *res,
                int row_number,
                int column_number);

This function returns 1 if the field is null and 0 if it contains a non-null value. (Note that PQgetvalue will return an empty string, not a null pointer, for a null field.)

PQgetlength

Returns the actual length of a field value in bytes. Row and column numbers start at 0.

int PQgetlength(const PGresult *res,
                int row_number,
                int column_number);

This is the actual data length for the particular data value, that is, the size of the object pointed to by PQgetvalue. For text data format this is the same as strlen(). For binary format this is essential information. Note that one should not rely on PQfsize to obtain the actual data length.

PQnparams

Returns the number of parameters of a prepared statement.

int PQnparams(const PGresult *res);

This function is only useful when inspecting the result of PQdescribePrepared. For other types of queries it will return zero.
Chapter 32. libpq - C Library

PQparamtype

Returns the data type of the indicated statement parameter. Parameter numbers start at 0.

\[
\text{Oid PQparamtype(const PGresult } \ast \text{res, int param_number);}
\]

This function is only useful when inspecting the result of PQdescribePrepared. For other types of queries it will return zero.

PQprint

Prints out all the rows and, optionally, the column names to the specified output stream.

\[
\text{void PQprint(FILE } \ast \text{fout, /* output stream */}
\text{const PGresult } \ast \text{res,}
\text{const PQprintOpt } \ast \text{po};}
\]

This function was formerly used by psql to print query results, but this is no longer the case. Note that it assumes all the data is in text format.

32.3.3. Retrieving Other Result Information

These functions are used to extract other information from PGresult objects.

PQcmdStatus

Returns the command status tag from the SQL command that generated the PGresult.

\[
\text{char } \ast \text{PQcmdStatus(PGresult } \ast \text{res);}
\]

Commonly this is just the name of the command, but it might include additional data such as the number of rows processed. The caller should not free the result directly. It will be freed when the associated PGresult handle is passed to PQclear.

PQcmdTuples

Returns the number of rows affected by the SQL command.

\[
\text{char } \ast \text{PQcmdTuples(PGresult } \ast \text{res);}
\]

This function returns a string containing the number of rows affected by the SQL statement that generated the PGresult. This function can only be used following the execution of a SELECT, CREATE TABLE AS, INSERT, UPDATE, DELETE, MOVE, FETCH, or COPY statement, or an EXECUTE of a prepared query that contains an INSERT, UPDATE, or DELETE statement. If the command that generated the PGresult was anything else, PQcmdTuples returns an empty string. The caller should not free the return value directly. It will be freed when the associated PGresult handle is passed to PQclear.
PQoidValue

Returns the OID of the inserted row, if the SQL command was an INSERT that inserted exactly one row into a table that has OIDs, or a EXECUTE of a prepared query containing a suitable INSERT statement. Otherwise, this function returns InvalidOid. This function will also return InvalidOid if the table affected by the INSERT statement does not contain OIDs.

Oid PQoidValue(const PGresult *res);

PQoidStatus

This function is deprecated in favor of PQoidValue and is not thread-safe. It returns a string with the OID of the inserted row, while PQoidValue returns the OID value.

char *PQoidStatus(const PGresult *res);

### 32.3.4. Escaping Strings for Inclusion in SQL Commands

PQescapeLiteral

char *PQescapeLiteral(PGconn *conn, const char *str, size_t length);

PQescapeLiteral escapes a string for use within an SQL command. This is useful when inserting data values as literal constants in SQL commands. Certain characters (such as quotes and backslashes) must be escaped to prevent them from being interpreted specially by the SQL parser. PQescapeLiteral performs this operation.

PQescapeLiteral returns an escaped version of the str parameter in memory allocated with malloc(). This memory should be freed using PQfreemem() when the result is no longer needed. A terminating zero byte is not required, and should not be counted in length. (If a terminating zero byte is found before length bytes are processed, PQescapeLiteral stops at the zero; the behavior is thus rather like strncpy().) The return string has all special characters replaced so that they can be properly processed by the PostgreSQL string literal parser. A terminating zero byte is also added. The single quotes that must surround PostgreSQL string literals are included in the result string.

On error, PQescapeLiteral returns NULL and a suitable message is stored in the conn object.

Tip: It is especially important to do proper escaping when handling strings that were received from an untrustworthy source. Otherwise there is a security risk: you are vulnerable to “SQL injection” attacks wherein unwanted SQL commands are fed to your database.

Note that it is neither necessary nor correct to do escaping when a data value is passed as a separate parameter in PQexecParams or its sibling routines.

PQescapeIdentifier

char *PQescapeIdentifier(PGconn *conn, const char *str, size_t length);

PQescapeIdentifier escapes a string for use as an SQL identifier, such as a table, column, or function name. This is useful when a user-supplied identifier might contain special characters that would otherwise not be interpreted as part of the identifier by the SQL parser, or when the identifier might contain upper case characters whose case should be preserved.

PQescapeIdentifier returns a version of the str parameter escaped as an SQL identifier in memory allocated with malloc(). This memory must be freed using PQfreemem() when the result is no longer needed. A terminating zero byte is not required, and should not be
counted in length. (If a terminating zero byte is found before length bytes are processed, PQescapeIdentifier stops at the zero; the behavior is thus rather like strncpy.) The return string has all special characters replaced so that it will be properly processed as an SQL identifier. A terminating zero byte is also added. The return string will also be surrounded by double quotes.

On error, PQescapeIdentifier returns NULL and a suitable message is stored in the conn object.

**Tip:** As with string literals, to prevent SQL injection attacks, SQL identifiers must be escaped when they are received from an untrustworthy source.

PQescapeStringConn

```c
size_t PQescapeStringConn(PGconn *conn, char *to, const char *from, size_t length, int *error);
```

PQescapeStringConn escapes string literals, much like PQescapeLiteral. Unlike PQescapeLiteral, the caller is responsible for providing an appropriately sized buffer. Furthermore, PQescapeStringConn does not generate the single quotes that must surround PostgreSQL string literals; they should be provided in the SQL command that the result is inserted into. The parameter from points to the first character of the string that is to be escaped, and the length parameter gives the number of bytes in this string. A terminating zero byte is not required, and should not be counted in length. (If a terminating zero byte is found before length bytes are processed, PQescapeStringConn stops at the zero; the behavior is thus rather like strncpy.) to shall point to a buffer that is able to hold at least one more byte than twice the value of length, otherwise the behavior is undefined. Behavior is likewise undefined if the to and from strings overlap.

If the error parameter is not NULL, then *error is set to zero on success, nonzero on error. Presently the only possible error conditions involve invalid multibyte encoding in the source string. The output string is still generated on error, but it can be expected that the server will reject it as malformed. On error, a suitable message is stored in the conn object, whether or not error is NULL.

PQescapeStringConn returns the number of bytes written to to, not including the terminating zero byte.

PQescapeString

```c
size_t PQescapeString (char *to, const char *from, size_t length);
```

The only difference from PQescapeStringConn is that PQescapeString does not take PGconn or error parameters. Because of this, it cannot adjust its behavior depending on the connection properties (such as character encoding) and therefore it might give the wrong results. Also, it has no way to report error conditions.

PQescapeString can be used safely in client programs that work with only one PostgreSQL connection at a time (in this case it can find out what it needs to know “behind the scenes”). In other contexts it is a security hazard and should be avoided in favor of PQescapeStringConn.

PQescapeByteaConn

Escapes binary data for use within an SQL command with the type bytea. As with PQescapeStringConn, this is only used when inserting data directly into an SQL command.
string.

```c
unsigned char *PQescapeByteaConn(PGconn *conn,
                                    const unsigned char *from,
                                    size_t from_length,
                                    size_t *to_length);
```

Certain byte values must be escaped when used as part of a byte literal in an SQL statement. `PQescapeByteaConn` escapes bytes using either hex encoding or backslash escaping. See Section 8.4 for more information.

The `from` parameter points to the first byte of the string that is to be escaped, and the `from_length` parameter gives the number of bytes in this binary string. (A terminating zero byte is neither necessary nor counted.) The `to_length` parameter points to a variable that will hold the resultant escaped string length. This result string length includes the terminating zero byte of the result.

`PQescapeByteaConn` returns an escaped version of the `from` parameter binary string in memory allocated with `malloc()`. This memory should be freed using `PQfreemem()` when the result is no longer needed. The return string has all special characters replaced so that they can be properly processed by the PostgreSQL string literal parser, and the bytea input function. A terminating zero byte is also added. The single quotes that must surround PostgreSQL string literals are not part of the result string.

On error, a null pointer is returned, and a suitable error message is stored in the `conn` object. Currently, the only possible error is insufficient memory for the result string.

`PQescapeBytea` is an older, deprecated version of `PQescapeByteaConn`.

```c
unsigned char *PQescapeBytea(const unsigned char *from,
                                  size_t from_length,
                                  size_t *to_length);
```

The only difference from `PQescapeByteaConn` is that `PQescapeBytea` does not take a `PGconn` parameter. Because of this, `PQescapeBytea` can only be used safely in client programs that use a single PostgreSQL connection at a time (in this case it can find out what it needs to know "behind the scenes"). It might give the wrong results if used in programs that use multiple database connections (use `PQescapeByteaConn` in such cases).

`PQunescapeBytea` Converts a string representation of binary data into binary data — the reverse of `PQescapeBytea`. This is needed when retrieving bytea data in text format, but not when retrieving it in binary format.

```c
unsigned char *PQunescapeBytea(const unsigned char *from, size_t *to_length);
```

The `from` parameter points to a string such as might be returned by `PQgetvalue` when applied to a bytea column. `PQunescapeBytea` converts this string representation into its binary representation. It returns a pointer to a buffer allocated with `malloc()`, or `NULL` on error, and puts the size of the buffer in `to_length`. The result must be freed using `PQfreemem` when it is no longer needed.

This conversion is not exactly the inverse of `PQescapeBytea`, because the string is not expected to be “escaped” when received from `PQgetvalue`. In particular this means there is no need for string quoting considerations, and so no need for a `PGconn` parameter.
32.4. Asynchronous Command Processing

The `PQexec` function is adequate for submitting commands in normal, synchronous applications. It has a few deficiencies, however, that can be of importance to some users:

- `PQexec` waits for the command to be completed. The application might have other work to do (such as maintaining a user interface), in which case it won’t want to block waiting for the response.
- Since the execution of the client application is suspended while it waits for the result, it is hard for the application to decide that it would like to try to cancel the ongoing command. (It can be done from a signal handler, but not otherwise.)
- `PQexec` can return only one `PGresult` structure. If the submitted command string contains multiple SQL commands, all but the last `PGresult` are discarded by `PQexec`.
- `PQexec` always collects the command’s entire result, buffering it in a single `PGresult`. While this simplifies error-handling logic for the application, it can be impractical for results containing many rows.

Applications that do not like these limitations can instead use the underlying functions that `PQexec` is built from: `PQsendQuery` and `PQgetResult`. There are also `PQsendQueryParams`, `PQsendPrepare`, `PQsendQueryPrepared`, `PQsendDescribePrepared`, and `PQsendDescribePortal`, which can be used with `PQgetResult` to duplicate the functionality of `PQexecParams`, `PQprepare`, `PQexecPrepared`, `PQdescribePrepared`, and `PQdescribePortal` respectively.

`PQsendQuery`

Submits a command to the server without waiting for the result(s). 1 is returned if the command was successfully dispatched and 0 if not (in which case, use `PQerrorMessage` to get more information about the failure).

```c
int PQsendQuery(PGconn *conn, const char *command);
```

After successfully calling `PQsendQuery`, call `PQgetResult` one or more times to obtain the results. `PQsendQuery` cannot be called again (on the same connection) until `PQgetResult` has returned a null pointer, indicating that the command is done.

`PQsendQueryParams`

Submits a command and separate parameters to the server without waiting for the result(s).

```c
int PQsendQueryParams(PGconn *conn, const char *command, int nParams, const Oid *paramTypes, const char * const *paramValues, const int *paramLengths, const int *paramFormats, int resultFormat);
```

This is equivalent to `PQsendQuery` except that query parameters can be specified separately from the query string. The function’s parameters are handled identically to `PQexecParams`. Like `PQexecParams`, it will not work on 2.0-protocol connections, and it allows only one command in the query string.
Chapter 32. libpq - C Library

PQsendPrepare

Sends a request to create a prepared statement with the given parameters, without waiting for completion.

```c
int PQsendPrepare(PGconn *conn,
    const char *stmtName,
    const char *query,
    int nParams,
    const Oid *paramTypes);
```

This is an asynchronous version of PQprepare: it returns 1 if it was able to dispatch the request, and 0 if not. After a successful call, call PQgetResult to determine whether the server successfully created the prepared statement. The function’s parameters are handled identically to PQprepare. Like PQprepare, it will not work on 2.0-protocol connections.

PQsendQueryPrepared

Sends a request to execute a prepared statement with given parameters, without waiting for the result(s).

```c
int PQsendQueryPrepared(PGconn *conn,
    const char *stmtName,
    int nParams,
    const char * const *paramValues,
    const int *paramLengths,
    const int *paramFormats,
    int resultFormat);
```

This is similar to PQsendQueryParams, but the command to be executed is specified by naming a previously-prepared statement, instead of giving a query string. The function’s parameters are handled identically to PQexecPrepared. Like PQexecPrepared, it will not work on 2.0-protocol connections.

PQsendDescribePrepared

Submits a request to obtain information about the specified prepared statement, without waiting for completion.

```c
int PQsendDescribePrepared(PGconn *conn, const char *stmtName);
```

This is an asynchronous version of PQdescribePrepared: it returns 1 if it was able to dispatch the request, and 0 if not. After a successful call, call PQgetResult to obtain the results. The function’s parameters are handled identically to PQdescribePrepared. Like PQdescribePrepared, it will not work on 2.0-protocol connections.

PQsendDescribePortal

Submits a request to obtain information about the specified portal, without waiting for completion.

```c
int PQsendDescribePortal(PGconn *conn, const char *portalName);
```

This is an asynchronous version of PQdescribePortal: it returns 1 if it was able to dispatch the request, and 0 if not. After a successful call, call PQgetResult to obtain the results. The function’s parameters are handled identically to PQdescribePortal. Like PQdescribePortal, it will not work on 2.0-protocol connections.

PQgetResult

Waits for the next result from a prior PQsendQuery, PQsendQueryParams, PQsendPrepare, PQsendQueryPrepared, PQsendDescribePrepared, or PQsendDescribePortal call, and returns it. A null pointer is returned when the command is complete and there will be no more results.

```c
PGresult *PQgetResult(PGconn *conn);
```
PQgetResult must be called repeatedly until it returns a null pointer, indicating that the command is done. (If called when no command is active, PQgetResult will just return a null pointer at once.) Each non-null result from PQgetResult should be processed using the same PGresult accessor functions previously described. Don’t forget to free each result object with PQclear when done with it. Note that PQgetResult will block only if a command is active and the necessary response data has not yet been read by PQconsumeInput.

**Note:** Even when PQresultStatus indicates a fatal error, PQgetResult should be called until it returns a null pointer, to allow libpq to process the error information completely.

Using PQsendQuery and PQgetResult solves one of PQexec’s problems: If a command string contains multiple SQL commands, the results of those commands can be obtained individually. (This allows a simple form of overlapped processing, by the way: the client can be handling the results of one command while the server is still working on later queries in the same command string.)

Another frequently-desired feature that can be obtained with PQsendQuery and PQgetResult is retrieving large query results a row at a time. This is discussed in Section 32.5.

By itself, calling PQgetResult will still cause the client to block until the server completes the next SQL command. This can be avoided by proper use of two more functions:

**PQconsumeInput**

If input is available from the server, consume it.

```c
int PQconsumeInput(PGconn *conn);
```

PQconsumeInput normally returns 1 indicating “no error”, but returns 0 if there was some kind of trouble (in which case PQerrorMessage can be consulted). Note that the result does not say whether any input data was actually collected. After calling PQconsumeInput, the application can check PQisBusy and/or PQnotifies to see if their state has changed.

PQconsumeInput can be called even if the application is not prepared to deal with a result or notification just yet. The function will read available data and save it in a buffer, thereby causing a select() read-ready indication to go away. The application can thus use PQconsumeInput to clear the select() condition immediately, and then examine the results at leisure.

**PQisBusy**

Returns 1 if a command is busy, that is, PQgetResult would block waiting for input. A 0 return indicates that PQgetResult can be called with assurance of not blocking.

```c
int PQisBusy(PGconn *conn);
```

PQisBusy will not itself attempt to read data from the server; therefore PQconsumeInput must be invoked first, or the busy state will never end.

A typical application using these functions will have a main loop that uses select() or poll() to wait for all the conditions that it must respond to. One of the conditions will be input available from the server, which in terms of select() means readable data on the file descriptor identified by PQsocket. When the main loop detects input ready, it should call PQconsumeInput to read the input. It can then call PQisBusy, followed by PQgetResult if PQisBusy returns false (0). It can also call PQnotifies to detect NOTIFY messages (see Section 32.8).
A client that uses PQsendQuery/PQgetResult can also attempt to cancel a command that is still being processed by the server; see Section 32.6. But regardless of the return value of PQcancel, the application must continue with the normal result-reading sequence using PQgetResult. A successful cancellation will simply cause the command to terminate sooner than it would have otherwise.

By using the functions described above, it is possible to avoid blocking while waiting for input from the database server. However, it is still possible that the application will block waiting to send output to the server. This is relatively uncommon but can happen if very long SQL commands or data values are sent. (It is much more probable if the application sends data via COPY IN, however.) To prevent this possibility and achieve completely nonblocking database operation, the following additional functions can be used.

**PQsetnonblocking**

Sets the nonblocking status of the connection.

```c
int PQsetnonblocking(PGconn *conn, int arg);
```

Sets the state of the connection to nonblocking if `arg` is 1, or blocking if `arg` is 0. Returns 0 if OK, -1 if error.

In the nonblocking state, calls to PQsendQuery, PQputline, PQputnbytes, PQputCopyData, and PQendcopy will not block but instead return an error if they need to be called again.

Note that PQexec does not honor nonblocking mode; if it is called, it will act in blocking fashion anyway.

**PQisnonblocking**

Returns the blocking status of the database connection.

```c
int PQisnonblocking(const PGconn *conn);
```

Returns 1 if the connection is set to nonblocking mode and 0 if blocking.

**PQflush**

Attempts to flush any queued output data to the server. Returns 0 if successful (or if the send queue is empty), -1 if it failed for some reason, or 1 if it was unable to send all the data in the send queue yet (this case can only occur if the connection is nonblocking).

```c
int PQflush(PGconn *conn);
```

After sending any command or data on a nonblocking connection, call PQflush. If it returns 1, wait for the socket to become read- or write-ready. If it becomes write-ready, call PQflush again. If it becomes read-ready, call PQconsumeInput, then call PQflush again. Repeat until PQflush returns 0. (It is necessary to check for read-ready and drain the input with PQconsumeInput, because the server can block trying to send us data, e.g. NOTICE messages, and won’t read our data until we read its.) Once PQflush returns 0, wait for the socket to be read-ready and then read the response as described above.

### 32.5. Retrieving Query Results Row-By-Row

Ordinarily, libpq collects a SQL command’s entire result and returns it to the application as a single PGresult. This can be unworkable for commands that return a large number of rows. For such cases, applications can use PQsendQuery and PQgetResult in single-row mode. In this mode, the result row(s) are returned to the application one at a time, as they are received from the server.
To enter single-row mode, call PQsetSingleRowMode immediately after a successful call of PQsendQuery (or a sibling function). This mode selection is effective only for the currently executing query. Then call PQgetResult repeatedly, until it returns null, as documented in Section 32.4. If the query returns any rows, they are returned as individual PGresult objects, which look like normal query results except for having status code PGRES_SINGLE_TUPLE instead of PGRES_TUPLES_OK. After the last row, or immediately if the query returns zero rows, a zero-row object with status PGRES_TUPLES_OK is returned; this is the signal that no more rows will arrive. (But note that it is still necessary to continue calling PQgetResult until it returns null.) All of these PGresult objects will contain the same row description data (column names, types, etc) that an ordinary PGresult object for the query would have. Each object should be freed with PQclear as usual.

PQsetSingleRowMode

Select single-row mode for the currently-executing query.

```c
int PQsetSingleRowMode(PGconn *conn);
```

This function can only be called immediately after PQsendQuery or one of its sibling functions, before any other operation on the connection such as PQconsumeInput or PQgetResult. If called at the correct time, the function activates single-row mode for the current query and returns 1. Otherwise the mode stays unchanged and the function returns 0. In any case, the mode reverts to normal after completion of the current query.

**Caution**

While processing a query, the server may return some rows and then encounter an error, causing the query to be aborted. Ordinarily, libpq discards any such rows and reports only the error. But in single-row mode, those rows will have already been returned to the application. Hence, the application will see some PGRES_SINGLE_TUPLE PGresult objects followed by a PGRES_FATAL_ERROR object. For proper transactional behavior, the application must be designed to discard or undo whatever has been done with the previously-processed rows, if the query ultimately fails.

### 32.6. Canceling Queries in Progress

A client application can request cancellation of a command that is still being processed by the server, using the functions described in this section.

PQgetCancel

Creates a data structure containing the information needed to cancel a command issued through a particular database connection.

```c
PGcancel *PQgetCancel(PGconn *conn);
```

PQgetCancel creates a PGcancel object given a PGconn connection object. It will return NULL if the given conn is NULL or an invalid connection. The PGcancel object is an opaque structure that is not meant to be accessed directly by the application; it can only be passed to PQcancel or PQfreeCancel.
PQfreeCancel

Frees a data structure created by PQgetCancel.

void PQfreeCancel(PGcancel *cancel);

PQfreeCancel frees a data object previously created by PQgetCancel.

PQcancel

Requests that the server abandon processing of the current command.

int PQcancel(PGcancel *cancel, char *errbuf, int errbufsize);

The return value is 1 if the cancel request was successfully dispatched and 0 if not. If not, errbuf is filled with an explanatory error message. errbuf must be a char array of size errbufsize (the recommended size is 256 bytes).

Successful dispatch is no guarantee that the request will have any effect, however. If the cancellation is effective, the current command will terminate early and return an error result. If the cancellation fails (say, because the server was already done processing the command), then there will be no visible result at all.

PQcancel can safely be invoked from a signal handler, if the errbuf is a local variable in the signal handler. The PGcancel object is read-only as far as PQcancel is concerned, so it can also be invoked from a thread that is separate from the one manipulating the PGconn object.

PQrequestCancel

PQrequestCancel is a deprecated variant of PQcancel.

int PQrequestCancel(PGconn *conn);

Requests that the server abandon processing of the current command. It operates directly on the PGconn object, and in case of failure stores the error message in the PGconn object (whence it can be retrieved by PQerrorMessage). Although the functionality is the same, this approach creates hazards for multiple-thread programs and signal handlers, since it is possible that overwriting the PGconn’s error message will mess up the operation currently in progress on the connection.

32.7. The Fast-Path Interface

PostgreSQL provides a fast-path interface to send simple function calls to the server.

Tip: This interface is somewhat obsolete, as one can achieve similar performance and greater functionality by setting up a prepared statement to define the function call. Then, executing the statement with binary transmission of parameters and results substitutes for a fast-path function call.

The function PQfn requests execution of a server function via the fast-path interface:

PGresult *PQfn(PGconn *conn,
   int fnid,
   int *result_buf,
   int *result_len,
   int result_is_int,
Chapter 32. libpq - C Library

const PQArgBlock *args,
int nargs);

typedef struct
{
    int len;
    int isint;
    union
    {
        int *ptr;
        int integer;
    } u;
} PQArgBlock;

The fnid argument is the OID of the function to be executed. args and nargs define the parameters to be passed to the function; they must match the declared function argument list. When the isint field of a parameter structure is true, the u.integer value is sent to the server as an integer of the indicated length (this must be 2 or 4 bytes); proper byte-swapping occurs. When isint is false, the indicated number of bytes at *u.ptr are sent with no processing; the data must be in the format expected by the server for binary transmission of the function’s argument data type. (The declaration of u.ptr as being of type int * is historical; it would be better to consider it void *.) result_buf points to the buffer in which to place the function's return value. The caller must have allocated sufficient space to store the return value. (There is no check!) The actual result length in bytes will be returned in the integer pointed to by result_len. If a 2- or 4-byte integer result is expected, set result_is_int to 1, otherwise set it to 0. Setting result_is_int to 1 causes libpq to byte-swap the value if necessary, so that it is delivered as a proper int value for the client machine; note that a 4-byte integer is delivered into *result_buf for either allowed result size. When result_is_int is 0, the binary-format byte string sent by the server is returned unmodified. (In this case it’s better to consider result_buf as being of type void *.)

PQfn always returns a valid PGresult pointer. The result status should be checked before the result is used. The caller is responsible for freeing the PGresult with PQclear when it is no longer needed.

Note that it is not possible to handle null arguments, null results, nor set-valued results when using this interface.

32.8. Asynchronous Notification

PostgreSQL offers asynchronous notification via the LISTEN and NOTIFY commands. A client session registers its interest in a particular notification channel with the LISTEN command (and can stop listening with the UNLISTEN command). All sessions listening on a particular channel will be notified asynchronously when a NOTIFY command with that channel name is executed by any session. A “payload” string can be passed to communicate additional data to the listeners.

libpq applications submit LISTEN, UNLISTEN, and NOTIFY commands as ordinary SQL commands. The arrival of NOTIFY messages can subsequently be detected by calling PQnotifies.

The function PQnotifies returns the next notification from a list of unhandled notification messages received from the server. It returns a null pointer if there are no pending notifications. Once a notification is returned from PQnotifies, it is considered handled and will be removed from the list of notifications.

PGnotify *PQnotifies(PGconn *conn);
typedef struct pgNotify
{
  char *relname;  /* notification channel name */
  int be_pid;    /* process ID of notifying server process */
  char *extra;   /* notification payload string */
} PGnotify;

After processing a PGnotify object returned by PQnotifies, be sure to free it with PQfreemem. It is sufficient to free the PGnotify pointer; the relname and extra fields do not represent separate allocations. (The names of these fields are historical; in particular, channel names need not have anything to do with relation names.)

Example 32-2 gives a sample program that illustrates the use of asynchronous notification.

PQnotifies does not actually read data from the server; it just returns messages previously absorbed by another libpq function. In ancient releases of libpq, the only way to ensure timely receipt of NOTIFY messages was to constantly submit commands, even empty ones, and then check PQnotifies after each PQexec. While this still works, it is deprecated as a waste of processing power.

A better way to check for NOTIFY messages when you have no useful commands to execute is to call PQconsumeInput, then check PQnotifies. You can use select() to wait for data to arrive from the server, thereby using no CPU power unless there is something to do. (See PQsocket to obtain the file descriptor number to use with select().) Note that this will work OK whether you submit commands with PQsendQuery/PQgetResult or simply use PQexec. You should, however, remember to check PQnotifies after each PQgetResult or PQexec, to see if any notifications came in during the processing of the command.

32.9. Functions Associated with the COPY Command

The COPY command in PostgreSQL has options to read from or write to the network connection used by libpq. The functions described in this section allow applications to take advantage of this capability by supplying or consuming copied data.

The overall process is that the application first issues the SQL COPY command via PQexec or one of the equivalent functions. The response to this (if there is no error in the command) will be a PGresult object bearing a status code of PGRES_COPY_OUT or PGRES_COPY_IN (depending on the specified copy direction). The application should then use the functions of this section to receive or transmit data rows. When the data transfer is complete, another PGresult object is returned to indicate success or failure of the transfer. Its status will be PGRES_COMMAND_OK for success or PGRES_FATAL_ERROR if some problem was encountered. At this point further SQL commands can be issued via PQexec. (It is not possible to execute other SQL commands using the same connection while the COPY operation is in progress.)

If a COPY command is issued via PQexec in a string that could contain additional commands, the application must continue fetching results via PQgetResult after completing the COPY sequence. Only when PQgetResult returns NULL is it certain that the PQexec command string is done and it is safe to issue more commands.

The functions of this section should be executed only after obtaining a result status of PGRES_COPY_OUT or PGRES_COPY_IN from PQexec or PQgetResult.

A PGresult object bearing one of these status values carries some additional data about the COPY operation that is starting. This additional data is available using functions that are also used in connection with query results:
Chapter 32. libpq - C Library

PQnfields
Returns the number of columns (fields) to be copied.

PQbinaryTuples
0 indicates the overall copy format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary. See COPY for more information.

PQfformat
Returns the format code (0 for text, 1 for binary) associated with each column of the copy operation. The per-column format codes will always be zero when the overall copy format is textual, but the binary format can support both text and binary columns. (However, as of the current implementation of COPY, only binary columns appear in a binary copy; so the per-column formats always match the overall format at present.)

Note: These additional data values are only available when using protocol 3.0. When using protocol 2.0, all these functions will return 0.

32.9.1. Functions for Sending COPY Data
These functions are used to send data during COPY FROM STDIN. They will fail if called when the connection is not in COPY_IN state.

PQputCopyData
Sends data to the server during COPY_IN state.

```c
int PQputCopyData(PGconn *conn, const char *buffer, int nbytes);
```

Transmits the COPY data in the specified `buffer`, of length `nbytes`, to the server. The result is 1 if the data was queued, zero if it was not queued because of full buffers (this will only happen in nonblocking mode), or -1 if an error occurred. (Use PQerrorMessage to retrieve details if the return value is -1. If the value is zero, wait for write-ready and try again.)

The application can divide the COPY data stream into buffer loads of any convenient size. Buffer-load boundaries have no semantic significance when sending. The contents of the data stream must match the data format expected by the COPY command; see COPY for details.

PQputCopyEnd
Sends end-of-data indication to the server during COPY_IN state.

```c
int PQputCopyEnd(PGconn *conn, const char *errmsg);
```

Ends the COPY_IN operation successfully if `errmsg` is NULL. If `errmsg` is not NULL then the COPY is forced to fail, with the string pointed to by `errmsg` used as the error message. (One should not assume that this exact error message will come back from the server, however, as the server might have already failed the COPY for its own reasons. Also note that the option to force failure does not work when using pre-3.0-protocol connections.)

The result is 1 if the termination message was sent; or in nonblocking mode, this may only indicate that the termination message was successfully queued. (In nonblocking mode, to be
certain that the data has been sent, you should next wait for write-ready and call PQflush, repeating until it returns zero.) Zero indicates that the function could not queue the termination message because of full buffers; this will only happen in nonblocking mode. (In this case, wait for write-ready and try the PQputCopyEnd call again.) If a hard error occurs, -1 is returned; you can use PQerrorMessage to retrieve details.

After successfully calling PQputCopyEnd, call PQgetResult to obtain the final result status of the COPY command. One can wait for this result to be available in the usual way. Then return to normal operation.

### 32.9.2. Functions for Receiving COPY Data

These functions are used to receive data during COPY TO STDOUT. They will fail if called when the connection is not in COPY_OUT state.

**PQgetCopyData**

Receives data from the server during COPY_OUT state.

```c
int PQgetCopyData(PGconn *conn,
                  char **buffer,
                  int async);
```

Attempts to obtain another row of data from the server during a COPY. Data is always returned one data row at a time; if only a partial row is available, it is not returned. Successful return of a data row involves allocating a chunk of memory to hold the data. The buffer parameter must be non-NULL. *buffer is set to point to the allocated memory, or to NULL in cases where no buffer is returned. A non-NULL result buffer should be freed using PQfreemem when no longer needed.

When a row is successfully returned, the return value is the number of data bytes in the row (this will always be greater than zero). The returned string is always null-terminated, though this is probably only useful for textual COPY. A result of zero indicates that the COPY is still in progress, but no row is yet available (this is only possible when async is true). A result of -1 indicates that the COPY is done. A result of -2 indicates that an error occurred (consult PQerrorMessage for the reason).

When async is true (not zero), PQgetCopyData will not block waiting for input; it will return zero if the COPY is still in progress but no complete row is available. (In this case wait for read-ready and then call PQconsumeInput before calling PQgetCopyData again.) When async is false (zero), PQgetCopyData will block until data is available or the operation completes.

After PQgetCopyData returns -1, call PQgetResult to obtain the final result status of the COPY command. One can wait for this result to be available in the usual way. Then return to normal operation.

### 32.9.3. Obsolete Functions for COPY

These functions represent older methods of handling COPY. Although they still work, they are deprecated due to poor error handling, inconvenient methods of detecting end-of-data, and lack of support for binary or nonblocking transfers.
Chapter 32. libpq - C Library

PQgetline

Reads a newline-terminated line of characters (transmitted by the server) into a buffer string of size length.

```c
int PQgetline(PGconn *conn,
             char *buffer,
             int length);
```

This function copies up to length-1 characters into the buffer and converts the terminating newline into a zero byte. PQgetline returns EOF at the end of input, 0 if the entire line has been read, and 1 if the buffer is full but the terminating newline has not yet been read.

Note that the application must check to see if a new line consists of the two characters ", which indicates that the server has finished sending the results of the COPY command. If the application might receive lines that are more than length-1 characters long, care is needed to be sure it recognizes the \ line correctly (and does not, for example, mistake the end of a long data line for a terminator line).

PQgetlineAsync

Reads a row of COPY data (transmitted by the server) into a buffer without blocking.

```c
int PQgetlineAsync(PGconn *conn,
                    char *buffer,
                    int bufsize);
```

This function is similar to PQgetline, but it can be used by applications that must read COPY data asynchronously, that is, without blocking. Having issued the COPY command and gotten a PGRES_COPY_OUT response, the application should call PQconsumeInput and PQgetlineAsync until the end-of-data signal is detected.

Unlike PQgetline, this function takes responsibility for detecting end-of-data.

On each call, PQgetlineAsync will return data if a complete data row is available in libpq’s input buffer. Otherwise, no data is returned until the rest of the row arrives. The function returns -1 if the end-of-copy-data marker has been recognized, or 0 if no data is available, or a positive number giving the number of bytes of data returned. If -1 is returned, the caller must next call PQendcopy, and then return to normal processing.

The data returned will not extend beyond a data-row boundary. If possible a whole row will be returned at one time. But if the buffer offered by the caller is too small to hold a row sent by the server, then a partial data row will be returned. With textual data this can be detected by testing whether the last returned byte is \n or not. (In a binary COPY, actual parsing of the COPY data format will be needed to make the equivalent determination.) The returned string is not null-terminated. (If you want to add a terminating null, be sure to pass a bufsize one smaller than the room actually available.)

PQputline

Sends a null-terminated string to the server. Returns 0 if OK and EOF if unable to send the string.

```c
int PQputline(PGconn *conn,
               const char *string);
```

The COPY data stream sent by a series of calls to PQputline has the same format as that returned by PQgetlineAsync, except that applications are not obliged to send exactly one data row per PQputline call; it is okay to send a partial line or multiple lines per call.

Note: Before PostgreSQL protocol 3.0, it was necessary for the application to explicitly send the two characters ", as a final line to indicate to the server that it had finished sending COPY data. While this still works, it is deprecated and the special meaning of \ can be expected to
be removed in a future release. It is sufficient to call PQendcopy after having sent the actual data.

PQputnbytes
Sends a non-null-terminated string to the server. Returns 0 if OK and EOF if unable to send the string.

```c
int PQputnbytes(PGconn *conn, const char *buffer, int nbytes);
```

This is exactly like PQputline, except that the data buffer need not be null-terminated since the number of bytes to send is specified directly. Use this procedure when sending binary data.

PQendcopy
Synchronizes with the server.

```c
int PQendcopy(PGconn *conn);
```

This function waits until the server has finished the copying. It should either be issued when the last string has been sent to the server using PQputline or when the last string has been received from the server using PQgetline. It must be issued or the server will get “out of sync” with the client. Upon return from this function, the server is ready to receive the next SQL command. The return value is 0 on successful completion, nonzero otherwise. (Use PQerrorMessage to retrieve details if the return value is nonzero.)

When using PQgetResult, the application should respond to a PGRES_COPY_OUT result by executing PQgetline repeatedly, followed by PQendcopy after the terminator line is seen. It should then return to the PQgetResult loop until PQgetResult returns a null pointer. Similarly a PGRES_COPY_IN result is processed by a series of PQputline calls followed by PQendcopy, then return to the PQgetResult loop. This arrangement will ensure that a COPY command embedded in a series of SQL commands will be executed correctly.

Older applications are likely to submit a COPY via PQexec and assume that the transaction is done after PQendcopy. This will work correctly only if the COPY is the only SQL command in the command string.

### 32.10. Control Functions

These functions control miscellaneous details of libpq’s behavior.

PQclientEncoding
Returns the client encoding.

```c
int PQclientEncoding(const PGconn *conn);
```

Note that it returns the encoding ID, not a symbolic string such as EUC_JP. If unsuccessful, it returns -1. To convert an encoding ID to an encoding name, you can use:

```c
char *pg_encoding_to_char(int encoding_id);
```

PQsetClientEncoding
Sets the client encoding.

```c
int PQsetClientEncoding(PGconn *conn, const char *encoding);
```
conn is a connection to the server, and encoding is the encoding you want to use. If the function successfully sets the encoding, it returns 0, otherwise -1. The current encoding for this connection can be determined by using PQclientEncoding.

PQsetErrorVerbosity
Determines the verbosity of messages returned by PQerrorMessage and PQresultErrorMessage.

typedef enum
{
    PQERRORS_TERSE,
    PQERRORS_DEFAULT,
    PQERRORS_VERBOSE
} PGVerbosity;

PGVerbosity PQsetErrorVerbosity(PGconn *conn, PGVerbosity verbosity);
PQsetErrorVerbosity sets the verbosity mode, returning the connection’s previous setting. In TERSE mode, returned messages include severity, primary text, and position only; this will normally fit on a single line. The default mode produces messages that include the above plus any detail, hint, or context fields (these might span multiple lines). The VERBOSE mode includes all available fields. Changing the verbosity does not affect the messages available from already-existing PGresult objects, only subsequently-created ones. (But see PQresultVerboseErrorMessage if you want to print a previous error with a different verbosity.)

PQsetErrorContextVisibility
Determines the handling of CONTEXT fields in messages returned by PQerrorMessage and PQresultErrorMessage.

typedef enum
{
    PQSHOW_CONTEXT_NEVER,
    PQSHOW_CONTEXT_ERRORS,
    PQSHOW_CONTEXT_ALWAYS
} PGContextVisibility;

PGContextVisibility PQsetErrorContextVisibility(PGconn *conn, PGContextVisibility show_context);
PQsetErrorContextVisibility sets the context display mode, returning the connection’s previous setting. This mode controls whether the CONTEXT field is included in messages (unless the verbosity setting is TERSE, in which case CONTEXT is never shown). The NEVER mode never includes CONTEXT, while ALWAYS always includes it if available. In ERRORS mode (the default), CONTEXT fields are included only for error messages, not for notices and warnings. Changing this mode does not affect the messages available from already-existing PGresult objects, only subsequently-created ones. (But see PQresultVerboseErrorMessage if you want to print a previous error with a different display mode.)

PQtrace
Enables tracing of the client/server communication to a debugging file stream.

void PQtrace(PGconn *conn, FILE *stream);

Note: On Windows, if the libpq library and an application are compiled with different flags, this function call will crash the application because the internal representation of the FILE pointers differ. Specifically, multithreaded/single-threaded, release/debug, and static/dynamic flags should be the same for the library and all applications using that library.
Chapter 32. libpq - C Library

PQuntrace

Disables tracing started by PQtrace.

void PQuntrace(PGconn *conn);

32.11. Miscellaneous Functions

As always, there are some functions that just don’t fit anywhere.

PQfreemem

Frees memory allocated by libpq.

void PQfreemem(void *ptr);

Frees memory allocated by libpq, particularly PQescapeByteaConn, PQescapeBytea, PQunescapeBytea, and PQnotifies. It is particularly important that this function, rather than free(), be used on Microsoft Windows. This is because allocating memory in a DLL and releasing it in the application works only if multithreaded/single-threaded, release/debug, and static/dynamic flags are the same for the DLL and the application. On non-Microsoft Windows platforms, this function is the same as the standard library function free().

PQconninfoFree

Frees the data structures allocated by PQconndefaults or PQconninfoParse.

void PQconninfoFree(PQconninfoOption *connOptions);

A simple PQfreemem will not do for this, since the array contains references to subsidiary strings.

PQencryptPassword

Prepares the encrypted form of a PostgreSQL password.

char * PQencryptPassword(const char *passwd, const char *user);

This function is intended to be used by client applications that wish to send commands like ALTER USER joe PASSWORD 'pwd'. It is good practice not to send the original cleartext password in such a command, because it might be exposed in command logs, activity displays, and so on. Instead, use this function to convert the password to encrypted form before it is sent. The arguments are the cleartext password, and the SQL name of the user it is for. The return value is a string allocated by malloc, or NULL if out of memory. The caller can assume the string doesn’t contain any special characters that would require escaping. Use PQfreemem to free the result when done with it.

PQmakeEmptyPGresult

Constructs an empty PGresult object with the given status.

PGresult *PQmakeEmptyPGresult(PGconn *conn, ExecStatusType status);

This is libpq’s internal function to allocate and initialize an empty PGresult object. This function returns NULL if memory could not be allocated. It is exported because some applications find it useful to generate result objects (particularly objects with error status) themselves. If conn is not null and status indicates an error, the current error message of the specified connection is copied into the PGresult. Also, if conn is not null, any event procedures registered in the connection are copied into the PGresult. (They do not get PGEVT_RESULTCREATE calls, but see PQfireResultCreateEvents.) Note that PQclear should eventually be called on the object, just as with a PGresult returned by libpq itself.
PQfireResultCreateEvents

Fires a PGEVT_RESULTCREATE event (see Section 32.13) for each event procedure registered in the PGresult object. Returns non-zero for success, zero if any event procedure fails.

```c
int PQfireResultCreateEvents(PGconn *conn, PGresult *res);
```

The `conn` argument is passed through to event procedures but not used directly. It can be NULL if the event procedures won’t use it.

Event procedures that have already received a PGEVT_RESULTCREATE or PGEVT_RESULTCOPY event for this object are not fired again.

The main reason that this function is separate from PQmakeEmptyPGresult is that it is often appropriate to create a PGresult and fill it with data before invoking the event procedures.

PQcopyResult

Makes a copy of a PGresult object. The copy is not linked to the source result in any way and PQclear must be called when the copy is no longer needed. If the function fails, NULL is returned.

```c
PGresult *PQcopyResult(const PGresult *src, int flags);
```

This is not intended to make an exact copy. The returned result is always put into PGRES_TUPLES_OK status, and does not copy any error message in the source. (It does copy the command status string, however.) The `flags` argument determines what else is copied. It is a bitwise OR of several flags. PG_COPYRES_ATTRS specifies copying the source result’s attributes (column definitions). PG_COPYRES_TUPLES specifies copying the source result’s tuples. (This implies copying the attributes, too.) PG_COPYRES_NOTICEHOOKS specifies copying the source result’s notify hooks. PG_COPYRES_EVENTS specifies copying the source result’s events. (But any instance data associated with the source is not copied.)

PQsetResultAttrs

Sets the attributes of a PGresult object.

```c
int PQsetResultAttrs(PGresult *res, int numAttributes, PGresAttDesc *attDescs);
```

The provided `attDescs` are copied into the result. If the `attDescs` pointer is NULL or `numAttributes` is less than one, the request is ignored and the function succeeds. If `res` already contains attributes, the function will fail. If the function fails, the return value is zero. If the function succeeds, the return value is non-zero.

PQsetvalue

Sets a tuple field value of a PGresult object.

```c
int PQsetvalue(PGresult *res, int tup_num, int field_num, char *value, int len);
```

The function will automatically grow the result’s internal tuples array as needed. However, the `tup_num` argument must be less than or equal to PQntuples, meaning this function can only grow the tuples array one tuple at a time. But any field of any existing tuple can be modified in any order. If a value at `field_num` already exists, it will be overwritten. If `len` is -1 or `value` is NULL, the field value will be set to an SQL null value. The `value` is copied into the result’s private storage, thus is no longer needed after the function returns. If the function fails, the return value is zero. If the function succeeds, the return value is non-zero.

PQresultAlloc

Allocate subsidiary storage for a PGresult object.

```c
void *PQresultAlloc(PGresult *res, size_t nBytes);
```
Any memory allocated with this function will be freed when `res` is cleared. If the function fails, the return value is `NULL`. The result is guaranteed to be adequately aligned for any type of data, just as for `malloc`.

**PQlibVersion**

Return the version of libpq that is being used.

```c
int PQlibVersion(void);
```

The result of this function can be used to determine, at run time, if specific functionality is available in the currently loaded version of libpq. The function can be used, for example, to determine which connection options are available for `PQconnectdb` or if the `hex bytea` output added in PostgreSQL 9.0 is supported.

The number is formed by converting the major, minor, and revision numbers into two-decimal-digit numbers and appending them together. For example, version 9.1 will be returned as 90100, and version 9.1.2 will be returned as 90102 (leading zeroes are not shown).

**Note:** This function appeared in PostgreSQL version 9.1, so it cannot be used to detect required functionality in earlier versions, since linking to it will create a link dependency on version 9.1.

### 32.12. Notice Processing

Notice and warning messages generated by the server are not returned by the query execution functions, since they do not imply failure of the query. Instead they are passed to a notice handling function, and execution continues normally after the handler returns. The default notice handling function prints the message on `stderr`, but the application can override this behavior by supplying its own handling function.

For historical reasons, there are two levels of notice handling, called the notice receiver and notice processor. The default behavior is for the notice receiver to format the notice and pass a string to the notice processor for printing. However, an application that chooses to provide its own notice receiver will typically ignore the notice processor layer and just do all the work in the notice receiver.

The function `PQsetNoticeReceiver` sets or examines the current notice receiver for a connection object. Similarly, `PQsetNoticeProcessor` sets or examines the current notice processor.

```c
typedef void (*PQnoticeReceiver) (void *arg, const PGresult *res);
PQnoticeReceiver
PQsetNoticeReceiver(PGconn *conn,
                     PQnoticeReceiver proc,
                     void *arg);

typedef void (*PQnoticeProcessor) (void *arg, const char *message);
PQnoticeProcessor
PQsetNoticeProcessor(PGconn *conn,
                     PQnoticeProcessor proc,
                     void *arg);
```
Chapter 32. libpq - C Library

Each of these functions returns the previous notice receiver or processor function pointer, and sets the new value. If you supply a null function pointer, no action is taken, but the current pointer is returned.

When a notice or warning message is received from the server, or generated internally by libpq, the notice receiver function is called. It is passed the message in the form of a \texttt{PGRES\_NONFATAL\_ERROR} \texttt{PGresult}. (This allows the receiver to extract individual fields using \texttt{PQresultErrorField}, or obtain a complete preformatted message using \texttt{PQresultErrorMessage} or \texttt{PQresultVerboseErrorMessage}.) The same void pointer passed to \texttt{PQsetNoticeReceiver} is also passed. (This pointer can be used to access application-specific state if needed.)

The default notice receiver simply extracts the message (using \texttt{PQresultErrorMessage}) and passes it to the notice processor.

The notice processor is responsible for handling a notice or warning message given in text form. It is passed the string text of the message (including a trailing newline), plus a void pointer that is the same one passed to \texttt{PQsetNoticeProcessor}. (This pointer can be used to access application-specific state if needed.)

The default notice processor is simply:

```c
static void
defaultNoticeProcessor(void *arg, const char *message)
{
    fprintf(stderr, "\%s", message);
}
```

Once you have set a notice receiver or processor, you should expect that that function could be called as long as either the \texttt{PGconn} object or \texttt{PGresult} objects made from it exist. At creation of a \texttt{PGresult}, the \texttt{PGconn}'s current notice handling pointers are copied into the \texttt{PGresult} for possible use by functions like \texttt{PQgetValue}.

32.13. Event System

libpq’s event system is designed to notify registered event handlers about interesting libpq events, such as the creation or destruction of \texttt{PGconn} and \texttt{PGresult} objects. A principal use case is that this allows applications to associate their own data with a \texttt{PGconn} or \texttt{PGresult} and ensure that that data is freed at an appropriate time.

Each registered event handler is associated with two pieces of data, known to libpq only as opaque \texttt{void *} pointers. There is a \texttt{passthrough} pointer that is provided by the application when the event handler is registered with a \texttt{PGconn}. The passthrough pointer never changes for the life of the \texttt{PGconn} and all \texttt{PGresult}s generated from it; so if used, it must point to long-lived data. In addition there is an \texttt{instance data} pointer, which starts out \texttt{NULL} in every \texttt{PGconn} and \texttt{PGresult}. This pointer can be manipulated using the \texttt{PQinstanceData}, \texttt{PQsetInstanceData}, \texttt{PQresultInstanceData} and \texttt{PQsetResultInstanceData} functions. Note that unlike the passthrough pointer, instance data of a \texttt{PGconn} is not automatically inherited by \texttt{PGresult}s created from it. libpq does not know what passthrough and instance data pointers point to (if anything) and will never attempt to free them — that is the responsibility of the event handler.
32.13.1. Event Types

The enum **PGEventId** names the types of events handled by the event system. All its values have names beginning with **PGEVT**. For each event type, there is a corresponding event info structure that carries the parameters passed to the event handlers. The event types are:

- **PGEVT_REGISTER**
  - The register event occurs when **PQregisterEventProc** is called. It is the ideal time to initialize any instanceData an event procedure may need. Only one register event will be fired per event handler per connection. If the event procedure fails, the registration is aborted.

```c
typedef struct
    {
        PGconn *conn;
    } PGEventRegister;
```

When a **PGEVT_REGISTER** event is received, the evtInfo pointer should be cast to a **PGEventRegister**. This structure contains a PGconn that should be in the **CONNECTION_OK** status; guaranteed if one calls **PQregisterEventProc** right after obtaining a good PGconn. When returning a failure code, all cleanup must be performed as no **PGEVT_CONNDESTROY** event will be sent.

- **PGEVT_CONNRESET**
  - The connection reset event is fired on completion of **PQreset** or **PQresetPoll**. In both cases, the event is only fired if the reset was successful. If the event procedure fails, the entire connection reset will fail; the PGconn is put into **CONNECTION_BAD** status and **PQresetPoll** will return **PGRES_POLLING_FAILED**.

```c
typedef struct
    {
        PGconn *conn;
    } PGEventConnReset;
```

When a **PGEVT_CONNRESET** event is received, the evtInfo pointer should be cast to a **PGEventConnReset**. Although the contained PGconn was just reset, all event data remains unchanged. This event should be used to reset/reload/requery any associated instanceData. Note that even if the event procedure fails to process **PGEVT_CONNRESET**, it will still receive a **PGEVT_CONNDESTROY** event when the connection is closed.

- **PGEVT_CONNDESTROY**
  - The connection destroy event is fired in response to **PQfinish**. It is the event procedure’s responsibility to properly clean up its event data as libpq has no ability to manage this memory. Failure to clean up will lead to memory leaks.

```c
typedef struct
    {
        PGconn *conn;
    } PGEventConnDestroy;
```

When a **PGEVT_CONNDESTROY** event is received, the evtInfo pointer should be cast to a **PGEventConnDestroy**. This event is fired prior to **PQfinish** performing any other cleanup. The return value of the event procedure is ignored since there is no way of indicating a failure from **PQfinish**. Also, an event procedure failure should not abort the process of cleaning up unwanted memory.
PGEVT_RESULTCREATE

The result creation event is fired in response to any query execution function that generates a result, including PQgetResult. This event will only be fired after the result has been created successfully.

typedef struct
{
    PGconn *conn;
    PGresult *result;
} PGEventResultCreate;

When a PGEVT_RESULTCREATE event is received, the evtInfo pointer should be cast to a PGEventResultCreate *. The conn is the connection used to generate the result. This is the ideal place to initialize any instanceData that needs to be associated with the result. If the event procedure fails, the result will be cleared and the failure will be propagated. The event procedure must not try to PQclear the result object for itself. When returning a failure code, all cleanup must be performed as no PGEVT_RESULTDESTROY event will be sent.

PGEVT_RESULTCOPY

The result copy event is fired in response to PQcopyResult. This event will only be fired after the copy is complete. Only event procedures that have successfully handled the PGEVT_RESULTCREATE or PGEVT_RESULTCOPY event for the source result will receive PGEVT_RESULTCOPY events.

typedef struct
{
    const PGresult *src;
    PGresult *dest;
} PGEventResultCopy;

When a PGEVT_RESULTCOPY event is received, the evtInfo pointer should be cast to a PGEventResultCopy *. The src result is what was copied while the dest result is the copy destination. This event can be used to provide a deep copy of instanceData, since PQcopyResult cannot do that. If the event procedure fails, the entire copy operation will fail and the dest result will be cleared. When returning a failure code, all cleanup must be performed as no PGEVT_RESULTDESTROY event will be sent for the destination result.

PGEVT_RESULTDESTROY

The result destroy event is fired in response to a PQclear. It is the event procedure’s responsibility to properly clean up its event data as libpq has no ability to manage this memory. Failure to clean up will lead to memory leaks.

typedef struct
{
    PGresult *result;
} PGEventResultDestroy;

When a PGEVT_RESULTDESTROY event is received, the evtInfo pointer should be cast to a PGEventResultDestroy *. This event is fired prior to PQclear performing any other cleanup. The return value of the event procedure is ignored since there is no way of indicating a failure from PQclear. Also, an event procedure failure should not abort the process of cleaning up unwanted memory.
32.13.2. Event Callback Procedure

**PGEventProc**

`PGEventProc` is a typedef for a pointer to an event procedure, that is, the user callback function that receives events from `libpq`. The signature of an event procedure must be:

```c
int eventproc(PGEventId evtId, void *evtInfo, void *passThrough)
```

The `evtId` parameter indicates which `PGEVT` event occurred. The `evtInfo` pointer must be cast to the appropriate structure type to obtain further information about the event. The `passThrough` parameter is the pointer provided to `PQregisterEventProc` when the event procedure was registered. The function should return a non-zero value if it succeeds and zero if it fails.

A particular event procedure can be registered only once in any `PGconn`. This is because the address of the procedure is used as a lookup key to identify the associated instance data.

**Caution**

On Windows, functions can have two different addresses: one visible from outside a DLL and another visible from inside the DLL. One should be careful that only one of these addresses is used with `libpq`'s event-procedure functions, else confusion will result. The simplest rule for writing code that will work is to ensure that event procedures are declared `static`. If the procedure's address must be available outside its own source file, expose a separate function to return the address.

32.13.3. Event Support Functions

**PQregisterEventProc**

Registers an event callback procedure with `libpq`.

```c
int PQregisterEventProc(PGconn *conn, PGEventProc proc,
    const char *name, void *passThrough);
```

An event procedure must be registered once on each `PGconn` you want to receive events about. There is no limit, other than memory, on the number of event procedures that can be registered with a connection. The function returns a non-zero value if it succeeds and zero if it fails.

The `proc` argument will be called when a `libpq` event is fired. Its memory address is also used to lookup `instanceData`. The `name` argument is used to refer to the event procedure in error messages. This value cannot be `NULL` or a zero-length string. The `name` string is copied into the `PGconn`, so what is passed need not be long-lived. The `passThrough` pointer is passed to the `proc` whenever an event occurs. This argument can be `NULL`.

**PQsetInstanceData**

Sets the connection `conn`'s `instanceData` for procedure `proc` to `data`. This returns non-zero for success and zero for failure. (Failure is only possible if `proc` has not been properly registered in `conn`.)

```c
int PQsetInstanceData(PGconn *conn, PGEventProc proc, void *data);
```
Chapter 32. libpq - C Library

PQinstanceData

Returns the connection conn’s instanceData associated with procedure proc, or NULL if there is none.

```c
void *PQinstanceData(const PGconn *conn, PGEventProc proc);
```

PQresultSetInstanceData

Sets the result’s instanceData for proc to data. This returns non-zero for success and zero for failure. (Failure is only possible if proc has not been properly registered in the result.)

```c
int PQresultSetInstanceData(PGresult *res, PGEventProc proc, void *data);
```

PQresultInstanceData

Returns the result’s instanceData associated with proc, or NULL if there is none.

```c
void *PQresultInstanceData(const PGresult *res, PGEventProc proc);
```

32.13.4. Event Example

Here is a skeleton example of managing private data associated with libpq connections and results.

```c
#include <libpq-events.h>

typedef struct
{
    int n;
    char *str;
} mydata;

static int myEventProc(PGEventId evtId, void *evtInfo, void *passThrough);

int main(void)
{
    mydata *data;
    PGresult *res;
    PGconn *conn = PQconnectdb("dbname=postgres options=-csearch_path=");

    if (PQstatus(conn) != CONNECTION_OK)
    {
        fprintf(stderr, "Connection to database failed: %s",
                 PQerrorMessage(conn));
        PQfinish(conn);
        return 1;
    }

    /* called once on any connection that should receive events.
       * Sends a PGEVT_REGISTER to myEventProc.
       */
    if (!PQregisterEventProc(conn, myEventProc, "mydata_proc", NULL))
    {
        fprintf(stderr, "Cannot register PGEventProc\n");
        PQfinish(conn);
    }
    
    if (PQstatus(conn) != CONNECTION_OK)
    {
        fprintf(stderr, "Connection to database failed: %s",
                 PQerrorMessage(conn));
        PQfinish(conn);
    }
```
/* conn instanceData is available */
data = PQinstanceData(conn, myEventProc);

/* Sends a PGEVT_RESULTCREATE to myEventProc */
res = PQexec(conn, "SELECT 1 + 1");

/* result instanceData is available */
data = PQresultInstanceData(res, myEventProc);

/* If PG_COPYRES_EVENTS is used, sends a PGEVT_RESULTCOPY to myEventProc */
res_copy = PQcopyResult(res, PG_COPYRES_TUPLES | PG_COPYRES_EVENTS);

/* result instanceData is available if PG_COPYRES_EVENTS was used during the PQcopyResult call. */
data = PQresultInstanceData(res_copy, myEventProc);

/* Both clears send a PGEVT_RESULTDESTROY to myEventProc */
PQclear(res);
PQclear(res_copy);

/* Sends a PGEVT_CONNDESTROY to myEventProc */
PQfinish(conn);

return 0;
}

static int
myEventProc(PGEventId evtId, void *evtInfo, void *passThrough)
{
    switch (evtId)
    {
    case PGEVT_REGISTER:
    {
        PGEventRegister *e = (PGEventRegister *)evtInfo;
        mydata *data = get_mydata(e->conn);

        /* associate app specific data with connection */
        PQsetInstanceData(e->conn, myEventProc, data);
        break;
    }

    case PGEVT_CONNRESET:
    {
        PGEventConnReset *e = (PGEventConnReset *)evtInfo;
        mydata *data = PQinstanceData(e->conn, myEventProc);

        if (data)
            memset(data, 0, sizeof(mydata));
        break;
    }

    case PGEVT_CONNDESTROY:
    {

    
}
Chapter 32. libpq - C Library

```c
PGEventConnDestroy *e = (PGEventConnDestroy *)evtInfo;
mydata *data = PQinstanceData(e->conn, myEventProc);

/* free instance data because the conn is being destroyed */
if (data)
    free_mydata(data);
break;
}

case PGEVT_RESULTCREATE:
{
    PGEventResultCreate *e = (PGEventResultCreate *)evtInfo;
    mydata *conn_data = PQinstanceData(e->conn, myEventProc);
    mydata *res_data = dup_mydata(conn_data);

    /* associate app specific data with result (copy it from conn) */
PQsetResultInstanceData(e->result, myEventProc, res_data);
break;
}

case PGEVT_RESULTCOPY:
{
    PGEventResultCopy *e = (PGEventResultCopy *)evtInfo;
    mydata *src_data = PQresultInstanceData(e->src, myEventProc);
    mydata *dest_data = dup_mydata(src_data);

    /* associate app specific data with result (copy it from a result) */
PQsetResultInstanceData(e->dest, myEventProc, dest_data);
break;
}

case PGEVT_RESULTDESTROY:
{
    PGEventResultDestroy *e = (PGEventResultDestroy *)evtInfo;
    mydata *data = PQresultInstanceData(e->result, myEventProc);

    /* free instance data because the result is being destroyed */
    if (data)
        free_mydata(data);
    break;
}

    /* unknown event ID, just return TRUE. */
default:
    break;
}

return TRUE; /* event processing succeeded */
```
32.14. Environment Variables

The following environment variables can be used to select default connection parameter values, which will be used by PQconnectdb, PQsetdbLogin, and PQsetdb if no value is directly specified by the calling code. These are useful to avoid hard-coding database connection information into simple client applications, for example.

- `PGHOST` behaves the same as the host connection parameter.
- `PGHOSTADDR` behaves the same as the hostaddr connection parameter. This can be set instead of or in addition to `PGHOST` to avoid DNS lookup overhead.
- `PGPORT` behaves the same as the port connection parameter.
- `PGDATABASE` behaves the same as the dbname connection parameter.
- `PGUSER` behaves the same as the user connection parameter.
- `PGPASSWORD` behaves the same as the password connection parameter. Use of this environment variable is not recommended for security reasons, as some operating systems allow non-root users to see process environment variables via `ps`; instead consider using the `~/.pgpass` file (see Section 32.15).
- `PGPASSFILE` specifies the name of the password file to use for lookups. If not set, it defaults to `~/.pgpass` (see Section 32.15).
- `PGSERVICE` behaves the same as the service connection parameter.
- `PGSERVICEFILE` specifies the name of the per-user connection service file. If not set, it defaults to `~/.pg_service.conf` (see Section 32.16).
- `PGOPTIONS` behaves the same as the options connection parameter.
- `PGAPPNAME` behaves the same as the application_name connection parameter.
- `PGSSLMODE` behaves the same as the sslmode connection parameter.
- `PGREQUIRESSL` behaves the same as the requiressl connection parameter. This environment variable is deprecated in favor of the `PGSSLMODE` variable; setting both variables suppresses the effect of this one.
- `PGSSLCOMPRESSION` behaves the same as the sslcompression connection parameter.
- `PGSSLCERT` behaves the same as the sslcert connection parameter.
- `PGSSLKEY` behaves the same as the sslkey connection parameter.
- `PGSSLROOTCERT` behaves the same as the sslrootcert connection parameter.
- `PGSLLCRL` behaves the same as the sslcrl connection parameter.
- `PGREQUIREPEER` behaves the same as the requirepeer connection parameter.
- `PGKRBSRVNAME` behaves the same as the krbsrvname connection parameter.
- `PGGSSLIB` behaves the same as the gsslib connection parameter.
- `PGCONNECT_TIMEOUT` behaves the same as the connect_timeout connection parameter.
- `PGCLIENTENCODING` behaves the same as the client_encoding connection parameter.
The following environment variables can be used to specify default behavior for each PostgreSQL session. (See also the ALTER ROLE and ALTER DATABASE commands for ways to set default behavior on a per-user or per-database basis.)

- **PGDATESTYLE** sets the default style of date/time representation. (Equivalent to `SET datestyle`)
- **PGTZ** sets the default time zone. (Equivalent to `SET timezone`)
- **PGGEQO** sets the default mode for the genetic query optimizer. (Equivalent to `SET geqo`)

Refer to the SQL command `SET` for information on correct values for these environment variables.

The following environment variables determine internal behavior of libpq; they override compiled-in defaults.

- **PGSYSCONFDIR** sets the directory containing the `pg_service.conf` file and in a future version possibly other system-wide configuration files.
- **PGLOCALEDIR** sets the directory containing the `locale` files for message localization.

### 32.15. The Password File

The file `.pgpass` in a user’s home directory or the file referenced by `PGPASSFILE` can contain passwords to be used if the connection requires a password (and no password has been specified otherwise). On Microsoft Windows the file is named `%APPDATA%\postgresql\pgpass.conf` (where `%APPDATA%` refers to the Application Data subdirectory in the user’s profile).

This file should contain lines of the following format:

```
hostname:port:database:username:password
```

(You can add a reminder comment to the file by copying the line above and preceding it with `#`.) Each of the first four fields can be a literal value, or `*`, which matches anything. The password field from the first line that matches the current connection parameters will be used. (Therefore, put more-specific entries first when you are using wildcards.) If an entry needs to contain `:` or `\`, escape this character with `\`. A host name of `localhost` matches both TCP (host name `localhost`) and Unix domain socket (`pghost` empty or the default socket directory) connections coming from the local machine. In a standby server, a database name of `replication` matches streaming replication connections made to the master server. The `database` field is of limited usefulness because users have the same password for all databases in the same cluster.

On Unix systems, the permissions on `.pgpass` must disallow any access to world or group; achieve this by the command `chmod 0600 ~/.pgpass`. If the permissions are less strict than this, the file will be ignored. On Microsoft Windows, it is assumed that the file is stored in a directory that is secure, so no special permissions check is made.

### 32.16. The Connection Service File

The connection service file allows `libpq` connection parameters to be associated with a single service name. That service name can then be specified by a `libpq` connection, and the associated settings will
be used. This allows connection parameters to be modified without requiring a recompile of the libpq application. The service name can also be specified using the PGSERVICE environment variable.

The connection service file can be a per-user service file at ~/.pg_service.conf or the location specified by the environment variable PGSERVICEFILE, or it can be a system-wide file at `pg_config --sysconfdir'/pg_service.conf or in the directory specified by the environment variable PGXSCONFDIR. If service definitions with the same name exist in the user and the system file, the user file takes precedence.

The file uses an “INI file” format where the section name is the service name and the parameters are connection parameters; see Section 32.1.2 for a list. For example:

```bash
# comment
[mydb]
host=somehost
port=5433
user=admin
```

An example file is provided at share/pg_service.conf.sample.

### 32.17. LDAP Lookup of Connection Parameters

If libpq has been compiled with LDAP support (option --with-ldap for configure) it is possible to retrieve connection options like host or dbname via LDAP from a central server. The advantage is that if the connection parameters for a database change, the connection information doesn’t have to be updated on all client machines.

LDAP connection parameter lookup uses the connection service file pg_service.conf (see Section 32.16). A line in a pg_service.conf stanza that starts with ldap:// will be recognized as an LDAP URL and an LDAP query will be performed. The result must be a list of keyword = value pairs which will be used to set connection options. The URL must conform to RFC 1959 and be of the form

```
ldap://[hostname[:port]]/search_base?attribute?search_scope?filter
```

where hostname defaults to localhost and port defaults to 389.

Processing of pg_service.conf is terminated after a successful LDAP lookup, but is continued if the LDAP server cannot be contacted. This is to provide a fallback with further LDAP URL lines that point to different LDAP servers, classical keyword = value pairs, or default connection options. If you would rather get an error message in this case, add a syntactically incorrect line after the LDAP URL.

A sample LDAP entry that has been created with the LDIF file

```ldif
version:1
dn:cn=mydatabase,dc=mycompany,dc=com
changetype:add
objectclass:top
objectclass:device
cn=mydatabase
description:host=dbserver.mycompany.com
description:port=5439
description:dbname=mydb
description:user=mydb_user
description:sslmode=require
```
might be queried with the following LDAP URL:

ldap://ldap.mycompany.com/dc=mycompany,dc=com?description=one?(cn=mydatabase)

You can also mix regular service file entries with LDAP lookups. A complete example for a stanza in pg_service.conf would be:

# only host and port are stored in LDAP, specify dbname and user explicitly [customerdb]
dbname=customer
user=appuser
ldap://ldap.acme.com/cn= dbs server, cn=hosts?pgconnectinfo?base?(objectclass=*)

32.18. SSL Support

PostgreSQL has native support for using SSL connections to encrypt client/server communications for increased security. See Section 18.9 for details about the server-side SSL functionality.

libpq reads the system-wide OpenSSL configuration file. By default, this file is named openssl.cnf and is located in the directory reported by openssl version -d. This default can be overridden by setting environment variable OPENSSL_CONF to the name of the desired configuration file.

32.18.1. Client Verification of Server Certificates

By default, PostgreSQL will not perform any verification of the server certificate. This means that it is possible to spoof the server identity (for example by modifying a DNS record or by taking over the server IP address) without the client knowing. In order to prevent spoofing, the client must be able to verify the server’s identity via a chain of trust. A chain of trust is established by placing a root (self-signed) certificate authority (CA) certificate on one computer and a leaf certificate signed by the root certificate on another computer. It is also possible to use an “intermediate” certificate which is signed by the root certificate and signs leaf certificates.

To allow the client to verify the identity of the server, place a root certificate on the client and a leaf certificate signed by the root certificate on the server. To allow the server to verify the identity of the client, place a root certificate on the server and a leaf certificate signed by the root certificate on the client. One or more intermediate certificates (usually stored with the leaf certificate) can also be used to link the leaf certificate to the root certificate.

Once a chain of trust has been established, there are two ways for the client to validate the leaf certificate sent by the server. If the parameter sslmode is set to verify-ca, libpq will verify that the server is trustworthy by checking the certificate chain up to the root certificate stored on the client. If sslmode is set to verify-full, libpq will also verify that the server host name matches the name stored in the server certificate. The SSL connection will fail if the server certificate cannot be verified. verify-full is recommended in most security-sensitive environments.

In verify-full mode, the host name is matched against the certificate’s Subject Alternative Name attribute(s), or against the Common Name attribute if no Subject Alternative Name of type dnsName is present. If the certificate’s name attribute starts with an asterisk (\*), the asterisk will be treated as a wildcard, which will match all characters except a dot (\.). This means the certificate will not match
subdomains. If the connection is made using an IP address instead of a host name, the IP address will be matched (without doing any DNS lookups).

To allow server certificate verification, one or more root certificates must be placed in the file ~/.postgresql/root.crt in the user’s home directory. (On Microsoft Windows the file is named %APPDATA%\postgresql\root.crt.) Intermediate certificates should also be added to the file if they are needed to link the certificate chain sent by the server to the root certificates stored on the client.

Certificate Revocation List (CRL) entries are also checked if the file ~/.postgresql/root.crl exists (%APPDATA%\postgresql\root.crt on Microsoft Windows).

The location of the root certificate file and the CRL can be changed by setting the connection parameters sslrootcert and sslcrl or the environment variables PGSSLROOTCERT and PGSSLCRL.

Note: For backwards compatibility with earlier versions of PostgreSQL, if a root CA file exists, the behavior of sslmode=require will be the same as that of verify-ca, meaning the server certificate is validated against the CA. Relying on this behavior is discouraged, and applications that need certificate validation should always use verify-ca or verify-full.

32.18.2. Client Certificates

If the server attempts to verify the identity of the client by requesting the client’s leaf certificate, libpq will send the certificates stored in file ~/.postgresql/postgresql.crt in the user’s home directory. The certificates must chain to the root certificate trusted by the server. A matching private key file ~/.postgresql/postgresql.key must also be present. The private key file must not allow any access to world or group; achieve this by the command chmod 0600 ~/.postgresql/postgresql.key. On Microsoft Windows these files are named %APPDATA%\postgresql\postgresql.crt and %APPDATA%\postgresql\postgresql.key, and there is no special permissions check since the directory is presumed secure. The location of the certificate and key files can be overridden by the connection parameters sslcert and sslkey or the environment variables PGSSLCERT and PGSSLKEY.

The first certificate in postgresql.crt must be the client’s certificate because it must match the client’s private key. “Intermediate” certificates can be optionally appended to the file — doing so avoids requiring storage of intermediate certificates on the server’s root.crt file.

For instructions on creating certificates, see Section 18.9.3.

32.18.3. Protection Provided in Different Modes

The different values for the sslmode parameter provide different levels of protection. SSL can provide protection against three types of attacks:

Eavesdropping

If a third party can examine the network traffic between the client and the server, it can read both connection information (including the user name and password) and the data that is passed. SSL uses encryption to prevent this.
Man in the middle (MITM)

If a third party can modify the data while passing between the client and server, it can pretend to be the server and therefore see and modify data even if it is encrypted. The third party can then forward the connection information and data to the original server, making it impossible to detect this attack. Common vectors to do this include DNS poisoning and address hijacking, whereby the client is directed to a different server than intended. There are also several other attack methods that can accomplish this. SSL uses certificate verification to prevent this, by authenticating the server to the client.

Impersonation

If a third party can pretend to be an authorized client, it can simply access data it should not have access to. Typically this can happen through insecure password management. SSL uses client certificates to prevent this, by making sure that only holders of valid certificates can access the server.

For a connection to be known secure, SSL usage must be configured on both the client and the server before the connection is made. If it is only configured on the server, the client may end up sending sensitive information (e.g. passwords) before it knows that the server requires high security. In libpq, secure connections can be ensured by setting the sslmode parameter to verify-full or verify-ca, and providing the system with a root certificate to verify against. This is analogous to using an https URL for encrypted web browsing.

Once the server has been authenticated, the client can pass sensitive data. This means that up until this point, the client does not need to know if certificates will be used for authentication, making it safe to specify that only in the server configuration.

All SSL options carry overhead in the form of encryption and key-exchange, so there is a trade-off that has to be made between performance and security. Table 32-1 illustrates the risks the different sslmode values protect against, and what statement they make about security and overhead.

Table 32-1. SSL Mode Descriptions

<table>
<thead>
<tr>
<th>sslmode</th>
<th>Eavesdropping protection</th>
<th>MITM protection</th>
<th>Statement</th>
</tr>
</thead>
<tbody>
<tr>
<td>disable</td>
<td>No</td>
<td>No</td>
<td>I don’t care about security, and I don’t want to pay the overhead of encryption.</td>
</tr>
<tr>
<td>allow</td>
<td>Maybe</td>
<td>No</td>
<td>I don’t care about security, but I will pay the overhead of encryption if the server insists on it.</td>
</tr>
<tr>
<td>prefer</td>
<td>Maybe</td>
<td>No</td>
<td>I don’t care about encryption, but I wish to pay the overhead of encryption if the server supports it.</td>
</tr>
</tbody>
</table>
sslmode | Eavesdropping protection | MITM protection | Statement |
--- | --- | --- | --- |
require | Yes | No | I want my data to be encrypted, and I accept the overhead. I trust that the network will make sure I always connect to the server I want. |
verify-ca | Yes | Depends on CA-policy | I want my data encrypted, and I accept the overhead. I want to be sure that I connect to a server that I trust. |
verify-full | Yes | Yes | I want my data encrypted, and I accept the overhead. I want to be sure that I connect to a server I trust, and that it’s the one I specify. |

The difference between verify-ca and verify-full depends on the policy of the root CA. If a public CA is used, verify-ca allows connections to a server that somebody else may have registered with the CA. In this case, verify-full should always be used. If a local CA is used, or even a self-signed certificate, using verify-ca often provides enough protection.

The default value for sslmode is prefer. As is shown in the table, this makes no sense from a security point of view, and it only promises performance overhead if possible. It is only provided as the default for backward compatibility, and is not recommended in secure deployments.

### 32.18.4. SSL Client File Usage

Table 32-2 summarizes the files that are relevant to the SSL setup on the client.

<table>
<thead>
<tr>
<th>File</th>
<th>Contents</th>
<th>Effect</th>
</tr>
</thead>
<tbody>
<tr>
<td>~/.postgresql/postgresql.crt</td>
<td>client certificate</td>
<td>requested by server</td>
</tr>
<tr>
<td>~/.postgresql/postgresql.key</td>
<td>client private key</td>
<td>proves client certificate sent by owner; does not indicate certificate owner is trustworthy</td>
</tr>
<tr>
<td>~/.postgresql/root.crt</td>
<td>trusted certificate authorities</td>
<td>checks that server certificate is signed by a trusted certificate authority</td>
</tr>
<tr>
<td>~/.postgresql/root.crl</td>
<td>certificates revoked by certificate authorities</td>
<td>server certificate must not be on this list</td>
</tr>
</tbody>
</table>
32.18.5. SSL Library Initialization

If your application initializes libssl and/or libcrypto libraries and libpq is built with SSL support, you should call PQinitOpenSSL to tell libpq that the libssl and/or libcrypto libraries have been initialized by your application, so that libpq will not also initialize those libraries. See http://h41379.www4.hpe.com/doc/83final/ba554_90007/ch04.html for details on the SSL API.

PQinitOpenSSL

Allows applications to select which security libraries to initialize.

```c
void PQinitOpenSSL(int do_ssl, int do_crypto);
```

When `do_ssl` is non-zero, libpq will initialize the OpenSSL library before first opening a database connection. When `do_crypto` is non-zero, the libcrypto library will be initialized. By default (if PQinitOpenSSL is not called), both libraries are initialized. When SSL support is not compiled in, this function is present but does nothing.

If your application uses and initializes either OpenSSL or its underlying libcrypto library, you must call this function with zeroes for the appropriate parameter(s) before first opening a database connection. Also be sure that you have done that initialization before opening a database connection.

PQinitSSL

Allows applications to select which security libraries to initialize.

```c
void PQinitSSL(int do_ssl);
```

This function is equivalent to PQinitOpenSSL(do_ssl, do_ssl). It is sufficient for applications that initialize both or neither of OpenSSL and libcrypto.

PQinitSSL has been present since PostgreSQL 8.0, while PQinitOpenSSL was added in PostgreSQL 8.4, so PQinitSSL might be preferable for applications that need to work with older versions of libpq.

32.19. Behavior in Threaded Programs

libpq is reentrant and thread-safe by default. You might need to use special compiler command-line options when you compile your application code. Refer to your system’s documentation for information about how to build thread-enabled applications, or look in `src/Makefile.global` for `PTHREAD_CFLAGS` and `PTHREAD_LIBS`. This function allows the querying of libpq’s thread-safe status:

PQisthreadsafe

```c
int PQisthreadsafe();
```

Returns 1 if the libpq is thread-safe and 0 if it is not.

One thread restriction is that no two threads attempt to manipulate the same PGconn object at the same time. In particular, you cannot issue concurrent commands from different threads through the same connection object. (If you need to run concurrent commands, use multiple connections.)
PGresult objects are normally read-only after creation, and so can be passed around freely between threads. However, if you use any of the PGresult-modifying functions described in Section 32.11 or Section 32.13, it’s up to you to avoid concurrent operations on the same PGresult, too.

The deprecated functions PQrequestCancel and PQoidStatus are not thread-safe and should not be used in multithread programs. PQrequestCancel can be replaced by PQcancel. PQoidStatus can be replaced by PQoidValue.

If you are using Kerberos inside your application (in addition to inside libpq), you will need to do locking around Kerberos calls because Kerberos functions are not thread-safe. See function PQregisterThreadLock in the libpq source code for a way to do cooperative locking between libpq and your application.

If you experience problems with threaded applications, run the program in src/tools/thread to see if your platform has thread-unsafe functions. This program is run by configure, but for binary distributions your library might not match the library used to build the binaries.

### 32.20. Building libpq Programs

To build (i.e., compile and link) a program using libpq you need to do all of the following things:

- Include the libpq-fe.h header file:

  \[
  \texttt{#include <libpq-fe.h>}
  \]

  If you failed to do that then you will normally get error messages from your compiler similar to:

  foo.c: In function ‘main’:
  foo.c:34: ‘PConn’ undeclared (first use in this function)
  foo.c:35: ‘PGresult’ undeclared (first use in this function)
  foo.c:54: ‘CONNECTION_BAD’ undeclared (first use in this function)
  foo.c:68: ‘PGRES_COMMAND_OK’ undeclared (first use in this function)
  foo.c:95: ‘PGRES_TUPLES_OK’ undeclared (first use in this function)

- Point your compiler to the directory where the PostgreSQL header files were installed, by supplying the `-I` directory option to your compiler. (In some cases the compiler will look into the directory in question by default, so you can omit this option.) For instance, your compile command line could look like:

  \[
  \texttt{cc -c -I/usr/local/pgsql/include testprog.c}
  \]

  If you are using makefiles then add the option to the CPPFLAGS variable:

  \[
  \texttt{CPPFLAGS += -I/usr/local/pgsql/include}
  \]

  If there is any chance that your program might be compiled by other users then you should not hardcode the directory location like that. Instead, you can run the utility `pg_config` to find out where the header files are on the local system:

  \[
  \$ \texttt{pg_config --includedir /usr/local/include}
  \]

  If you have `pkg-config` installed, you can run instead:

  \[
  \$ \texttt{pkg-config --cflags libpq -I/usr/local/include}
  \]

  Note that this will already include the `-I` in front of the path.

  Failure to specify the correct option to the compiler will result in an error message such as:

  \[
  \texttt{testlibpq.c:8:22: libpq-fe.h: No such file or directory}
  \]
When linking the final program, specify the option `-lpq` so that the libpq library gets pulled in, as well as the option `-L` to point the compiler to the directory where the libpq library resides. (Again, the compiler will search some directories by default.) For maximum portability, put the `-L` option before the `-lpq` option. For example:

```
c -o testprog testprog1.o testprog2.o -L/usr/local/pgsql/lib -lpq
```

You can find out the library directory using `pg_config` as well:

```
$ pg_config --libdir /usr/local/pgsql/lib
```

Or again use `pkg-config`:

```
$ pkg-config --libs libpq -L/usr/local/pgsql/lib -lpq
```

Note again that this prints the full options, not only the path.

Error messages that point to problems in this area could look like the following:

```
testlibpq.o: In function 'main':
testlibpq.o(.text+0x60): undefined reference to 'PQsetdbLogin'
testlibpq.o(.text+0x71): undefined reference to 'PQstatus'
testlibpq.o(.text+0x4a): undefined reference to 'PQerrorMessage'
```

This means you forgot `-lpq`.

```
/usr/bin/ld: cannot find -lpq
```

This means you forgot the `-L` option or did not specify the right directory.

## 32.21. Example Programs

These examples and others can be found in the directory `src/test/examples` in the source code distribution.

### Example 32-1. libpq Example Program 1

```c
/*
 * src/test/examples/testlibpq.c
 *
 * testlibpq.c
 *
 * Test the C version of libpq, the PostgreSQL frontend library.
 */
#include <stdio.h>
#include <stdlib.h>
#include "libpq-fe.h"

static void
exit_nicely(PGconn *conn)
{
    PQfinish(conn);
    exit(1);
}

int
main(int argc, char **argv)
```
Chapter 32. libpq - C Library

const char *conninfo;
PGconn *conn;
PGresult *res;
int nFields;
int i,
j;

/*
 * If the user supplies a parameter on the command line, use it as the
 * conninfo string; otherwise default to setting dbname=postgres and using
 * environment variables or defaults for all other connection parameters.
 */
if (argc > 1)
    conninfo = argv[1];
else
    conninfo = "dbname = postgres";

/* Make a connection to the database */
conn = PQconnectdb(conninfo);

/* Check to see that the backend connection was successfully made */
if (PQstatus(conn) != CONNECTION_OK)
{
    fprintf(stderr, "Connection to database failed: %s",
        PQerrorMessage(conn));
    exit_nicely(conn);
}

/* Set always-secure search path, so malicious users can’t take control. */
res = PQexec(conn,
    "SELECT pg_catalog.set_config('search_path', ", false)"
);
if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
    PQclear(res);
    exit_nicely(conn);
}

/* Should PQclear PGresult whenever it is no longer needed to avoid memory
 * leaks */
PQclear(res);

/*
 * Our test case here involves using a cursor, for which we must be inside
 * a transaction block. We could do the whole thing with a single
 * PQexec() of "select * from pg_database", but that’s too trivial to make
 * a good example.
 */

/* Start a transaction block */
res = PQexec(conn, "BEGIN");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
{
    fprintf(stderr, "BEGIN command failed: %s", PQerrorMessage(conn));
}
Chapter 32. libpq - C Library

PQclear(res);
exit_nicely(conn);
}
PQclear(res);

/*
 * Fetch rows from pg_database, the system catalog of databases
 */
res = PQexec(conn, "DECLARE myportal CURSOR FOR select * from pg_database");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
{
    fprintf(stderr, "DECLARE CURSOR failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

res = PQexec(conn, "FETCH ALL in myportal");
if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "FETCH ALL failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}

/* first, print out the attribute names */
nFields = PQnfields(res);
for (i = 0; i < nFields; i++)
    printf("%-15s", PQfname(res, i));
printf("\n\n");

/* next, print out the rows */
for (i = 0; i < PQntuples(res); i++)
{
    for (j = 0; j < nFields; j++)
        printf("%-15s", PQgetvalue(res, i, j));
    printf("\n");
}
PQclear(res);

/* close the portal ... we don’t bother to check for errors ... */
res = PQexec(conn, "CLOSE myportal");
PQclear(res);

/* end the transaction */
res = PQexec(conn, "END");
PQclear(res);

/* close the connection to the database and cleanup */
PQfinish(conn);
return 0;
}
Example 32-2. libpq Example Program 2

/*
 * src/test/examples/testlibpq2.c
 *
 * testlibpq2.c
 * Test of the asynchronous notification interface
 *
 * Start this program, then from psql in another window do
 * NOTIFY TBL2;
 * Repeat four times to get this program to exit.
 *
 * Or, if you want to get fancy, try this:
 * populate a database with the following commands
 * (provided in src/test/examples/testlibpq2.sql):
 *
 * CREATE SCHEMA TESTLIBPQ2;
 * SET search_path = TESTLIBPQ2;
 * CREATE TABLE TBL1 (i int4);
 * CREATE TABLE TBL2 (i int4);
 * CREATE RULE r1 AS ON INSERT TO TBL1 DO
 *   (INSERT INTO TBL2 VALUES (new.i); NOTIFY TBL2);
 *
 * Start this program, then from psql do this four times:
 *
 * INSERT INTO TESTLIBPQ2.TBL1 VALUES (10);
 */

#ifdef WIN32
#include <windows.h>
#endif
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <errno.h>
#include <sys/time.h>
#include <sys/types.h>
#ifdef HAVE_SYS_SELECT_H
#include <sys/select.h>
#endif
#include "libpq-fe.h"

static void
exit_nicely(PGconn *conn)
{
  PQfinish(conn);
  exit(1);
}

int
main(int argc, char **argv)
{
  const char *conninfo;
  PGconn  *conn;
  PGresult  *res;

PGnotify notify;
int nnotifies;

/*
 * If the user supplies a parameter on the command line, use it as the
 * conninfo string; otherwise default to setting dbname=postgres and using
 * environment variables or defaults for all other connection parameters.
 */
if (argc > 1)
    conninfo = argv[1];
else
    conninfo = "dbname = postgres";

/* Make a connection to the database */
conn = PQconnectdb(conninfo);

/* Check to see that the backend connection was successfully made */
if (PQstatus(conn) != CONNECTION_OK)
{
    fprintf(stderr, "Connection to database failed: %s",
            PQerrorMessage(conn));
    exit_nicely(conn);
}

/* Set always-secure search path, so malicious users can’t take control. */
res = PQexec(conn,
"SELECT pg_catalog.set_config(‘search_path’, "", false)"
);  
if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
    PQclear(res);
    exit_nicely(conn);
}

/*
 * Should PQclear PGresult whenever it is no longer needed to avoid memory
 * leaks
 */
PQclear(res);

/*
 * Issue LISTEN command to enable notifications from the rule’s NOTIFY.
 */
res = PQexec(conn, "LISTEN TBL2");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
{
    fprintf(stderr, "LISTEN command failed: %s", PQerrorMessage(conn));
    PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

/* Quit after four notifies are received. */
nnotifies = 0;
while (nnotifies < 4)
{
    /*
   */
```c
/* Sleep until something happens on the connection. We use select(2) to wait for input, but you could also use poll() or similar facilities. */
int sock;
fd_set input_mask;

sock = PQsocket(conn);
if (sock < 0)
    break; /* shouldn’t happen */

FD_ZERO(&input_mask);
FD_SET(sock, &input_mask);

if (select(sock + 1, &input_mask, NULL, NULL, NULL) < 0)
{
    fprintf(stderr, "select() failed: %s\n", strerror(errno));
    exit_nicely(conn);
}

/* Now check for input */
PQconsumeInput(conn);
while ((notify = PQnotifies(conn)) != NULL)
{
    fprintf(stderr,
            "ASYNC NOTIFY of ‘%s’ received from backend PID %d\n",
            notify->relname, notify->be_pid);
    PQfreemem(notify);
    nnotifies++;
    PQconsumeInput(conn);
}

fprintf(stderr, "Done.\n");

/* close the connection to the database and cleanup */
PQfinish(conn);
return 0;
}
```

Example 32-3. libpq Example Program 3

```bash
/*
* src/test/examples/testlibpq3.c
* 
* testlibpq3.c
* Test out-of-line parameters and binary I/O.
* 
* Before running this, populate a database with the following commands
* (provided in src/test/examples/testlibpq3.sql):
* 
* CREATE SCHEMA testlibpq3;
* SET search_path = testlibpq3;
* CREATE TABLE test1 (i int4, t text, b bytea);
```
* INSERT INTO test1 values (1, ’joe’s place’, ‘\000\001\002\003\004’);
* INSERT INTO test1 values (2, ’ho there’, ‘\004\003\002\001\000’);
*
* The expected output is:
*
* tuple 0: got
* i = (4 bytes) 1
* t = (11 bytes) ’joe’s place’
* b = (5 bytes) \000\001\002\003\004
*
* tuple 0: got
* i = (4 bytes) 2
* t = (8 bytes) ’ho there’
* b = (5 bytes) \004\003\002\001\000
*/

#ifdef WIN32
#include <windows.h>
#endif
#include <stdio.h>
#include <stdlib.h>
#include <stdint.h>
#include <string.h>
#include <sys/types.h>
#include "libpq-fe.h"
/* for ntohl/htonl */
#include <netinet/in.h>
#include <arpa/inet.h>

static void
exit_nicely(PGconn *conn)
{
    PQfinish(conn);
    exit(1);
}

/*/  
*
* This function prints a query result that is a binary-format fetch from  
* a table defined as in the comment above. We split it out because the  
* main() function uses it twice.
*/
static void
show_binary_results(PGresult *res)
{
    int i,
        j;
    int i_fnum,
        t_fnum,
        b_fnum;

    /* Use PQfnumber to avoid assumptions about field order in result */
    i_fnum = PQfnumber(res, "i");
    t_fnum = PQfnumber(res, "t");
    b_fnum = PQfnumber(res, "b");
for (i = 0; i < PQntuples(res); i++)
{
  char *iptr;
  char *tptr;
  char *bptr;
  int blen;
  int ival;

  /* Get the field values (we ignore possibility they are null!) */
  iptr = PQgetvalue(res, i, i_fnum);
  tptr = PQgetvalue(res, i, t_fnum);
  bptr = PQgetvalue(res, i, b_fnum);

  /* The binary representation of INT4 is in network byte order, which
   we'd better coerce to the local byte order. */
  ival = ntohl(*((uint32_t *) iptr));

  /* The binary representation of TEXT is, well, text, and since libpq
   was nice enough to append a zero byte to it, it'll work just fine
   as a C string. */
  /* The binary representation of BYTEA is a bunch of bytes, which could
   include embedded nulls so we have to pay attention to field length. */
  blen = PQgetlength(res, i, b_fnum);

  printf("tuple %d: got\n", i);
  printf("  i = (%d bytes) %d\n",
         PQgetlength(res, i, i_fnum), ival);
  printf("  t = (%d bytes) '\%s '\n",
         PQgetlength(res, i, t_fnum), tptr);
  printf("  b = (%d bytes) ", blen);
  for (j = 0; j < blen; j++)
    printf("\\\%03o", bptr[j]);
  printf("\n\n\n");
}

int
main(int argc, char **argv)
{
  const char *conninfo;
  PGconn *conn;
  PResult *res;
  const char *paramValues[1];
  int   paramLengths[1];
  int   paramFormats[1];
  uint32_t binaryIntVal;

  /* If the user supplies a parameter on the command line, use it as the
   conninfo string; otherwise default to setting dbname=postgres and using
   environment variables or defaults for all other connection parameters.
if (argc > 1)
    conninfo = argv[1];
else
    conninfo = "dbname = postgres";

/* Make a connection to the database */
conn = PQconnectdb(conninfo);

/* Check to see that the backend connection was successfully made */
if (PQstatus(conn) != CONNECTION_OK)
    {
        fprintf(stderr, "Connection to database failed: %s",
                PQerrorMessage(conn));
        exit_nicely(conn);
    }

/* Set always-secure search path, so malicious users can’t take control. */
res = PQexec(conn, "SET search_path = testlibpq3");
if (PQresultStatus(res) != PGRES_COMMAND_OK)
    {
        fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
        PQclear(res);
        exit_nicely(conn);
    }

PQclear(res);

/* The point of this program is to illustrate use of PQexecParams() with
 out-of-line parameters, as well as binary transmission of data.

 This first example transmits the parameters as text, but receives the
 results in binary format. By using out-of-line parameters we can avoid
 a lot of tedious mucking about with quoting and escaping, even though
 the data is text. Notice how we don’t have to do anything special with
 the quote mark in the parameter value.
*/

/* Here is our out-of-line parameter value */
paramValues[0] = "joe’s place";

res = PQexecParams(conn,
    "SELECT * FROM test1 WHERE t = $1",
    1,  /* one param */
    NULL,  /* let the backend deduce param type */
    paramValues,
    NULL,  /* don’t need param lengths since text */
    NULL,  /* default to all text params */
    1);  /* ask for binary results */

if (PQresultStatus(res) != PGRES_TUPLES_OK)
    {
        fprintf(stderr, "SELECT failed: %s", PQerrorMessage(conn));
        PQclear(res);
        exit_nicely(conn);
    }
show_binary_results(res);
PQclear(res);

/*
 * In this second example we transmit an integer parameter in binary form,
 * and again retrieve the results in binary form.
 *
 * Although we tell PQexecParams we are letting the backend deduce
 * parameter type, we really force the decision by casting the parameter
 * symbol in the query text. This is a good safety measure when sending
 * binary parameters.
 */

/* Convert integer value "2" to network byte order */
binaryIntVal = htonl((uint32_t) 2);

/* Set up parameter arrays for PQexecParams */
paramValues[0] = (char *) &binaryIntVal;
paramLengths[0] = sizeof(binaryIntVal);
paramFormats[0] = 1; /* binary */

res = PQexecParams(conn,
    "SELECT * FROM test1 WHERE i = $1::int4",
    1, /* one param */
    NULL, /* let the backend deduce param type */
    paramValues,
    paramLengths,
    paramFormats,
    1); /* ask for binary results */

if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SELECT failed: %s", PQerrorMessage(conn));
PQclear(res);
    exit_nicely(conn);
}

show_binary_results(res);
PQclear(res);

/* close the connection to the database and cleanup */
PQfinish(conn);

return 0;
Chapter 33. Large Objects

PostgreSQL has a large object facility, which provides stream-style access to user data that is stored in a special large-object structure. Streaming access is useful when working with data values that are too large to manipulate conveniently as a whole.

This chapter describes the implementation and the programming and query language interfaces to PostgreSQL large object data. We use the libpq C library for the examples in this chapter, but most programming interfaces native to PostgreSQL support equivalent functionality. Other interfaces might use the large object interface internally to provide generic support for large values. This is not described here.

33.1. Introduction

All large objects are stored in a single system table named pg_largeobject. Each large object also has an entry in the system table pg_largeobject_metadata. Large objects can be created, modified, and deleted using a read/write API that is similar to standard operations on files.

PostgreSQL also supports a storage system called “TOAST”, which automatically stores values larger than a single database page into a secondary storage area per table. This makes the large object facility partially obsolete. One remaining advantage of the large object facility is that it allows values up to 4 TB in size, whereas TOASTed fields can be at most 1 GB. Also, reading and updating portions of a large object can be done efficiently, while most operations on a TOASTed field will read or write the whole value as a unit.

33.2. Implementation Features

The large object implementation breaks large objects up into “chunks” and stores the chunks in rows in the database. A B-tree index guarantees fast searches for the correct chunk number when doing random access reads and writes.

The chunks stored for a large object do not have to be contiguous. For example, if an application opens a new large object, seeks to offset 1000000, and writes a few bytes there, this does not result in allocation of 1000000 bytes worth of storage; only of chunks covering the range of data bytes actually written. A read operation will, however, read out zeroes for any unallocated locations preceding the last existing chunk. This corresponds to the common behavior of “sparsely allocated” files in Unix file systems.

As of PostgreSQL 9.0, large objects have an owner and a set of access permissions, which can be managed using GRANT and REVOKE. SELECT privileges are required to read a large object, and UPDATE privileges are required to write or truncate it. Only the large object’s owner (or a database superuser) can delete, comment on, or change the owner of a large object. To adjust this behavior for compatibility with prior releases, see the lo_compat_privileges run-time parameter.

33.3. Client Interfaces

This section describes the facilities that PostgreSQL’s libpq client interface library provides for accessing large objects. The PostgreSQL large object interface is modeled after the Unix file-system interface, with analogues of open, read, write, lseek, etc.
All large object manipulation using these functions must take place within an SQL transaction block, since large object file descriptors are only valid for the duration of a transaction.

If an error occurs while executing any one of these functions, the function will return an otherwise-impossible value, typically 0 or -1. A message describing the error is stored in the connection object and can be retrieved with PQerrorMessage.

Client applications that use these functions should include the header file libpq/libpq-fs.h and link with the libpq library.

### 33.3.1. Creating a Large Object

The function

```c
Oid lo_creat(PGconn *conn, int mode);
```

creates a new large object. The return value is the OID that was assigned to the new large object, or InvalidOid (zero) on failure. `mode` is unused and ignored as of PostgreSQL 8.1; however, for backward compatibility with earlier releases it is best to set it to INV_READ, INV_WRITE, or INV_READ | INV_WRITE. (These symbolic constants are defined in the header file libpq/libpq-fs.h.)

An example:

```c
inv_oid = lo_creat(conn, INV_READ|INV_WRITE);
```

The function

```c
Oid lo_create(PGconn *conn, Oid lobjId);
```

also creates a new large object. The OID to be assigned can be specified by `lobjId`; if so, failure occurs if that OID is already in use for some large object. If `lobjId` is InvalidOid (zero) then `lo_create` assigns an unused OID (this is the same behavior as `lo_creat`). The return value is the OID that was assigned to the new large object, or InvalidOid (zero) on failure.

`lo_create` is new as of PostgreSQL 8.1; if this function is run against an older server version, it will fail and return InvalidOid.

An example:

```c
inv_oid = lo_create(conn, desired_oid);
```

### 33.3.2. Importing a Large Object

To import an operating system file as a large object, call

```c
Oid lo_import(PGconn *conn, const char *filename);
```

`filename` specifies the operating system name of the file to be imported as a large object. The return value is the OID that was assigned to the new large object, or InvalidOid (zero) on failure. Note that the file is read by the client interface library, not by the server; so it must exist in the client file system and be readable by the client application.

The function
Oid lo_import_with_oid(PGconn *conn, const char *filename, Oid lobjId);

also imports a new large object. The OID to be assigned can be specified by lobjId; if so, failure occurs if that OID is already in use for some large object. If lobjId is InvalidOid (zero) then lo_import_with_oid assigns an unused OID (this is the same behavior as lo_import). The return value is the OID that was assigned to the new large object, or InvalidOid (zero) on failure.

lo_import_with_oid is new as of PostgreSQL 8.4 and uses lo_create internally which is new in 8.1; if this function is run against 8.0 or before, it will fail and return InvalidOid.

### 33.3.3. Exporting a Large Object

To export a large object into an operating system file, call

```c
int lo_export(PGconn *conn, Oid lobjId, const char *filename);
```

The lobjId argument specifies the OID of the large object to export and the filename argument specifies the operating system name of the file. Note that the file is written by the client interface library, not by the server. Returns 1 on success, -1 on failure.

### 33.3.4. Opening an Existing Large Object

To open an existing large object for reading or writing, call

```c
int lo_open(PGconn *conn, Oid lobjId, int mode);
```

The lobjId argument specifies the OID of the large object to open. The mode bits control whether the object is opened for reading (INV_READ), writing (INV_WRITE), or both. (These symbolic constants are defined in the header file `libpq/libpq-fs.h`.) lo_open returns a (non-negative) large object descriptor for later use in lo_read, lo_write, lo_lseek, lo_lseek64, lo_tell, lo_tell64, lo_truncate, lo_truncate64, and lo_close. The descriptor is only valid for the duration of the current transaction. On failure, -1 is returned.

The server currently does not distinguish between modes INV_WRITE and INV_READ | INV_WRITE: you are allowed to read from the descriptor in either case. However there is a significant difference between these modes and INV_READ alone: with INV_READ you cannot write on the descriptor, and the data read from it will reflect the contents of the large object at the time of the transaction snapshot that was active when lo_open was executed, regardless of later writes by this or other transactions. Reading from a descriptor opened with INV_WRITE returns data that reflects all writes of other committed transactions as well as writes of the current transaction. This is similar to the behavior of REPEATABLE READ versus READ COMMITTED transaction modes for ordinary SQL SELECT commands.

An example:

```c
inv_fd = lo_open(conn, inv_oid, INV_READ|INV_WRITE);
```

### 33.3.5. Writing Data to a Large Object

The function

```c
int lo_write(PGconn *conn, int fd, const char *buf, size_t len);
```
Chapter 33. Large Objects

writes \texttt{len} bytes from \texttt{buf} (which must be of size \texttt{len}) to large object descriptor \texttt{fd}. The \texttt{fd} argument must have been returned by a previous \texttt{lo_open}. The number of bytes actually written is returned (in the current implementation, this will always equal \texttt{len} unless there is an error). In the event of an error, the return value is -1.

Although the \texttt{len} parameter is declared as \texttt{size_t}, this function will reject length values larger than \texttt{INT_MAX}. In practice, it’s best to transfer data in chunks of at most a few megabytes anyway.

### 33.3.6. Reading Data from a Large Object

The function

\begin{verbatim}
int lo_read(PGconn *conn, int fd, char *buf, size_t len);
\end{verbatim}

reads up to \texttt{len} bytes from large object descriptor \texttt{fd} into \texttt{buf} (which must be of size \texttt{len}). The \texttt{fd} argument must have been returned by a previous \texttt{lo_open}. The number of bytes actually read is returned; this will be less than \texttt{len} if the end of the large object is reached first. In the event of an error, the return value is -1.

Although the \texttt{len} parameter is declared as \texttt{size_t}, this function will reject length values larger than \texttt{INT_MAX}. In practice, it’s best to transfer data in chunks of at most a few megabytes anyway.

### 33.3.7. Seeking in a Large Object

To change the current read or write location associated with a large object descriptor, call

\begin{verbatim}
int lo_lseek(PGconn *conn, int fd, int offset, int whence);
\end{verbatim}

This function moves the current location pointer for the large object descriptor identified by \texttt{fd} to the new location specified by \texttt{offset}. The valid values for \texttt{whence} are \texttt{SEEK_SET} (seek from object start), \texttt{SEEK_CUR} (seek from current position), and \texttt{SEEK_END} (seek from object end). The return value is the new location pointer, or -1 on error.

When dealing with large objects that might exceed 2GB in size, instead use

\begin{verbatim}
pg_int64 lo_lseek64(PGconn *conn, int fd, pg_int64 offset, int whence);
\end{verbatim}

This function has the same behavior as \texttt{lo_lseek}, but it can accept an \texttt{offset} larger than 2GB and/or deliver a result larger than 2GB. Note that \texttt{lo_lseek} will fail if the new location pointer would be greater than 2GB.

\texttt{lo_lseek64} is new as of PostgreSQL 9.3. If this function is run against an older server version, it will fail and return -1.

### 33.3.8. Obtaining the Seek Position of a Large Object

To obtain the current read or write location of a large object descriptor, call

\begin{verbatim}
int lo_tell(PGconn *conn, int fd);
\end{verbatim}

If there is an error, the return value is -1.

When dealing with large objects that might exceed 2GB in size, instead use
33.3.9. Truncating a Large Object

To truncate a large object to a given length, call

```c
int lo_truncate(PGconn *conn, int fd, size_t len);
```

This function truncates the large object descriptor `fd` to length `len`. The `fd` argument must have been returned by a previous `lo_open`. If `len` is greater than the large object’s current length, the large object is extended to the specified length with null bytes (\'\0\'). On success, `lo_truncate` returns zero. On error, the return value is -1.

The read/write location associated with the descriptor `fd` is not changed.

Although the `len` parameter is declared as `size_t`, `lo_truncate` will reject length values larger than `INT_MAX`.

When dealing with large objects that might exceed 2GB in size, instead use

```c
int lo_truncate64(PGconn *conn, int fd, pg_int64 len);
```

This function has the same behavior as `lo_truncate`, but it can accept a `len` value exceeding 2GB.

`lo_truncate` is new as of PostgreSQL 8.3; if this function is run against an older server version, it will fail and return -1.

`lo_truncate64` is new as of PostgreSQL 9.3; if this function is run against an older server version, it will fail and return -1.

33.3.10. Closing a Large Object Descriptor

A large object descriptor can be closed by calling

```c
int lo_close(PGconn *conn, int fd);
```

where `fd` is a large object descriptor returned by `lo_open`. On success, `lo_close` returns zero. On error, the return value is -1.

Any large object descriptors that remain open at the end of a transaction will be closed automatically.

33.3.11. Removing a Large Object

To remove a large object from the database, call

```c
int lo_unlink(PGconn *conn, Oid lobjId);
```

The `lobjId` argument specifies the OID of the large object to remove. Returns 1 if successful, -1 on failure.
33.4. Server-side Functions

Server-side functions tailored for manipulating large objects from SQL are listed in Table 33-1.

Table 33-1. SQL-oriented Large Object Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>lo_from_bytea()</td>
<td>oid</td>
<td>Create a large object and store data there, returning its OID. Pass 0 to have the system choose an OID.</td>
<td>lo_from_bytea(0, '\xff\xff\xff\xff\xff00')</td>
<td>24528</td>
</tr>
<tr>
<td>lo_put()</td>
<td>void</td>
<td>Write data at the given offset.</td>
<td>lo_put(24528, 1, '\xa')(1)</td>
<td></td>
</tr>
<tr>
<td>lo_get()</td>
<td>bytea</td>
<td>Extract contents or a substring thereof.</td>
<td>lo_get(24528, 0, 3)</td>
<td>\xff\xaaf</td>
</tr>
</tbody>
</table>

There are additional server-side functions corresponding to each of the client-side functions described earlier; indeed, for the most part the client-side functions are simply interfaces to the equivalent server-side functions. The ones just as convenient to call via SQL commands are lo_create, lo_unlink, lo_import, and lo_export. Here are examples of their use:

```sql
CREATE TABLE image (name text, raster oid);
SELECT lo_creat(-1); -- returns OID of new, empty large object
SELECT lo_create(43213); -- attempts to create large object with OID 43213
SELECT lo_unlink(173454); -- deletes large object with OID 173454
INSERT INTO image (name, raster) VALUES ('beautiful image', lo_import('/etc/motd'));
INSERT INTO image (name, raster) VALUES ('beautiful image', lo_import('/etc/motd', 68583));
SELECT lo_export(image.raster, '/tmp/motd') FROM image WHERE name = 'beautiful image';
```

The server-side lo_import and lo_export functions behave considerably differently from their client-side analogs. These two functions read and write files in the server’s file system, using the permissions of the database’s owning user. Therefore, their use is restricted to superusers. In contrast, the client-side import and export functions read and write files in the client’s file system, using the permissions of the client program. The client-side functions do not require superuser privilege.
The functionality of \texttt{lo\_read} and \texttt{lo\_write} is also available via server-side calls, but the names of the server-side functions differ from the client side interfaces in that they do not contain underscores. You must call these functions as \texttt{loread} and \texttt{lowrite}.

### 33.5. Example Program

Example 33-1 is a sample program which shows how the large object interface in libpq can be used. Parts of the program are commented out but are left in the source for the reader’s benefit. This program can also be found in \texttt{src/test/examples/testlo.c} in the source distribution.

#### Example 33-1. Large Objects with libpq Example Program

```c
/*-------------------------------------------------------------------------
* testlo.c
* test using large objects with libpq
* Portions Copyright (c) 1996-2016, PostgreSQL Global Development Group
* Portions Copyright (c) 1994, Regents of the University of California
* 
* IDENTIFICATION
* src/test/examples/testlo.c
* *-------------------------------------------------------------------------
*/
#include <stdio.h>
#include <stdlib.h>
#include <sys/types.h>
#include <sys/stat.h>
#include <fcntl.h>
#include <unistd.h>
#include "libpq-fe.h"
#include "libpq/libpq-fs.h"
#define BUFSIZE 1024

/*
 * importFile -
 * import file "in_filename" into database as large object "lobjOid"
 */
static Oid
importFile(PGconn *conn, char *filename)
{
    Oid lobjId;
    int lobj_fd;
    char buf[BUFSIZE];
    int nbytes,
        tmp;
    int fd;
```
Chapter 33. Large Objects

/*
 * open the file to be read in
 */
fd = open(filename, O_RDONLY, 0666);
if (fd < 0)
{ /* error */
    fprintf(stderr, "cannot open unix file"\"%s\"\n", filename);
}

/*
 * create the large object
 */
lobjId = lo_creat(conn, INV_READ | INV_WRITE);
if (lobjId == 0)
    fprintf(stderr, "cannot create large object");
lobj_fd = lo_open(conn, lobjId, INV_WRITE);

/*
 * read in from the Unix file and write to the inversion file
 */
while ((nbytes = read(fd, buf, BUFSIZE)) > 0)
{
    tmp = lo_write(conn, lobj_fd, buf, nbytes);
    if (tmp < nbytes)
        fprintf(stderr, "error while reading \"%s\", filename);
}

close(fd);
lo_close(conn, lobj_fd);

return lobjId;
}

static void
pickout(PGconn *conn, Oid lobjId, int start, int len)
{
    int lobj_fd;
    char *buf;
    int nbytes;
    int nread;

    lobj_fd = lo_open(conn, lobjId, INV_READ);
    if (lobj_fd < 0)
        fprintf(stderr, "cannot open large object \%u", lobjId);

    lo_lseek(conn, lobj_fd, start, SEEK_SET);
    buf = malloc(len + 1);
    nread = 0;
    while (len - nread > 0)
    {
        nbytes = lo_read(conn, lobj_fd, buf, len - nread);
        buf[nbytes] = '\0';
        fprintf(stderr, ">>> %s", buf);
        nread += nbytes;
        if (nbytes <= 0)
Chapter 33. Large Objects

break;       /* no more data? */
}
free(buf);
fprintf(stderr, "\n");
lo_close(conn, lobj_fd);
}

static void
overwrite(PGconn *conn, Oid lobjId, int start, int len)
{
    int lobj_fd;
    char *buf;
    int nbytes;
    int nwritten;
    int i;

    lobj_fd = lo_open(conn, lobjId, INV_WRITE);
    if (lobj_fd < 0)
        fprintf(stderr, "cannot open large object %u", lobjId);

    lo_lseek(conn, lobj_fd, start, SEEK_SET);
    buf = malloc(len + 1);

    for (i = 0; i < len; i++)
        buf[i] = 'X';
    buf[i] = '\0';

    nwritten = 0;
    while (len - nwritten > 0)
    {
        nbytes = lo_write(conn, lobj_fd, buf + nwritten, len - nwritten);
        nwritten += nbytes;
        if (nbytes <= 0)
        {
            fprintf(stderr, "\nWRITE FAILED!\n")
            break;
        }
    }
    free(buf);
    fprintf(stderr, "\n");
    lo_close(conn, lobj_fd);
}

/*
 * exportFile -
 * export large object "lobjOid" to file "out_filename"
 *
 */
static void
exportFile(PGconn *conn, Oid lobjId, char *filename)
{
    int lobj_fd;
    char buf[BUFSIZE];
    int nbytes,
        tmp;
    int fd;

Chapter 33. Large Objects

/*
 * open the large object
 */
lobj_fd = lo_open(conn, lobjId, INV_READ);
if (lobj_fd < 0)
    fprintf(stderr, "cannot open large object %u", lobjId);

/*
 * open the file to be written to
 */
fd = open(filename, O_CREAT | O_WRONLY | O_TRUNC, 0666);
if (fd < 0)
    { /* error */
        fprintf(stderr, "cannot open unix file" "%s",
                filename);
    }

/*
 * read in from the inversion file and write to the Unix file
 */
while ((nbytes = lo_read(conn, lobj_fd, buf, BUFSIZE)) > 0)
    { tmp = write(fd, buf, nbytes);
        if (tmp < nbytes)
            { fprintf(stderr, "error while writing \"%s\",
                        filename);
            }
    }
lo_close(conn, lobj_fd);
close(fd);
return;
}

static void
exit_nicely(PGconn *conn)
{
PQfinish(conn);
exit(1);
}

int
main(int argc, char **argv)
{
    char *in_filename,
         *out_filename;
    char *database;
    Oid lobjoid;
    PGconn *conn;
    PGresult *res;

    if (argc != 4)
        { fprintf(stderr, "Usage: %s database_name in_filename out_filename\n",}

710
Chapter 33. Large Objects

```c
argv[0]);
exit(1);
}

database = argv[1];
in_filename = argv[2];
out_filename = argv[3];

/*
 * set up the connection
 */
conn = PQsetdb(NULL, NULL, NULL, NULL, database);

/* check to see that the backend connection was successfully made */
if (PQstatus(conn) != CONNECTION_OK)
{
    fprintf(stderr, "Connection to database failed: %s",
            PQerrorMessage(conn));
    exit_nicely(conn);
}

/* Set always-secure search path, so malicious users can't take control. */
res = PQexec(conn,
             "SELECT pg_catalog.set_config('search_path', ", false)"");
if (PQresultStatus(res) != PGRES_TUPLES_OK)
{
    fprintf(stderr, "SET failed: %s", PQerrorMessage(conn));
    PQclear(res);
    exit_nicely(conn);
}
PQclear(res);

res = PQexec(conn, "begin");
PQclear(res);
printf("importing file "%s" ...
", in_filename);
/* lobjOid = importFile(conn, in_filename); */
lobjOid = lo_import(conn, in_filename);
if (lobjOid == 0)
    fprintf(stderr, "%s", PQerrorMessage(conn));
else
    printf("\tas large object %u.\n", lobjOid);

printf("picking out bytes 1000-2000 of the large object\n");
pickout(conn, lobjOid, 1000, 1000);

printf("overwriting bytes 1000-2000 of the large object with X's\n");
overwrite(conn, lobjOid, 1000, 1000);

printf("exporting large object to file "%s" ...
", out_filename);
/*
exportFile(conn, lobjOid, out_filename);
*/
if (lo_export(conn, lobjOid, out_filename) < 0)
    fprintf(stderr, "%s", PQerrorMessage(conn));
}
res = PQexec(conn, "end");
PQclear(res);
```
PQfinish(conn);
return 0;
}
Chapter 34. ECPG - Embedded SQL in C

This chapter describes the embedded SQL package for PostgreSQL. It was written by Linus Tolke (<linus@epact.se>) and Michael Meskes (<meskes@postgresql.org>). Originally it was written to work with C. It also works with C++, but it does not recognize all C++ constructs yet.

This documentation is quite incomplete. But since this interface is standardized, additional information can be found in many resources about SQL.

34.1. The Concept

An embedded SQL program consists of code written in an ordinary programming language, in this case C, mixed with SQL commands in specially marked sections. To build the program, the source code (*.pgc) is first passed through the embedded SQL preprocessor, which converts it to an ordinary C program (*.c), and afterwards it can be processed by a C compiler. (For details about the compiling and linking see Section 34.10). Converted ECPG applications call functions in the libpq library through the embedded SQL library (ecpglib), and communicate with the PostgreSQL server using the normal frontend-backend protocol.

Embedded SQL has advantages over other methods for handling SQL commands from C code. First, it takes care of the tedious passing of information to and from variables in your C program. Second, the SQL code in the program is checked at build time for syntactical correctness. Third, embedded SQL in C is specified in the SQL standard and supported by many other SQL database systems. The PostgreSQL implementation is designed to match this standard as much as possible, and it is usually possible to port embedded SQL programs written for other SQL databases to PostgreSQL with relative ease.

As already stated, programs written for the embedded SQL interface are normal C programs with special code inserted to perform database-related actions. This special code always has the form:

EXEC SQL ...;

These statements syntactically take the place of a C statement. Depending on the particular statement, they can appear at the global level or within a function. Embedded SQL statements follow the case-sensitivity rules of normal SQL code, and not those of C. Also they allow nested C-style comments that are part of the SQL standard. The C part of the program, however, follows the C standard of not accepting nested comments.

The following sections explain all the embedded SQL statements.

34.2. Managing Database Connections

This section describes how to open, close, and switch database connections.

34.2.1. Connecting to the Database Server

One connects to a database using the following statement:

EXEC SQL CONNECT TO target [AS connection-name] [USER user-name];

The target can be specified in the following ways:
Chapter 34. ECPG - Embedded SQL in C

- `dbname[@hostname][:port]`
- `tcp:postgresql://hostname[:port] [/dbname][?options]`
- `unix:postgresql://hostname[:port] [/dbname][?options]`
- an SQL string literal containing one of the above forms
- a reference to a character variable containing one of the above forms (see examples)
- DEFAULT

If you specify the connection target literally (that is, not through a variable reference) and you don’t quote the value, then the case-insensitivity rules of normal SQL are applied. In that case you can also double-quote the individual parameters separately as needed. In practice, it is probably less error-prone to use a (single-quoted) string literal or a variable reference. The connection target DEFAULT initiates a connection to the default database under the default user name. No separate user name or connection name can be specified in that case.

There are also different ways to specify the user name:

- `username`
- `username/password`
- `username IDENTIFIED BY password`
- `username USING password`

As above, the parameters `username` and `password` can be an SQL identifier, an SQL string literal, or a reference to a character variable.

If the connection target includes any `options`, those consist of `keyword=value` specifications separated by ampersands (&). The allowed key words are the same ones recognized by libpq (see Section 32.1.2). Spaces are ignored before any `keyword` or `value`, though not within or after one. Note that there is no way to write `&` within a `value`.

The `connection-name` is used to handle multiple connections in one program. It can be omitted if a program uses only one connection. The most recently opened connection becomes the current connection, which is used by default when an SQL statement is to be executed (see later in this chapter).

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin each session by removing publicly-writable schemas from `search_path`. For example, add `options=-c search_path= to options, or issue EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); after connecting. This consideration is not specific to ECPG; it applies to every interface for executing arbitrary SQL commands.

Here are some examples of `CONNECT` statements:

`EXEC SQL CONNECT TO mydb@sql.mydomain.com;`

`EXEC SQL CONNECT TO unix:postgresql://sql.mydomain.com/mydb AS myconnection USER john;`

`EXEC SQL BEGIN DECLARE SECTION;
const char *target = "mydb@sql.mydomain.com";
const char *user = "john";
const char *passwd = "secret";
EXEC SQL END DECLARE SECTION;
...
EXEC SQL CONNECT TO :target USER :user USING :passwd;`
The last form makes use of the variant referred to above as character variable reference. You will see in later sections how C variables can be used in SQL statements when you prefix them with a colon. Be advised that the format of the connection target is not specified in the SQL standard. So if you want to develop portable applications, you might want to use something based on the last example above to encapsulate the connection target string somewhere.

### 34.2.2. Choosing a Connection

SQL statements in embedded SQL programs are by default executed on the current connection, that is, the most recently opened one. If an application needs to manage multiple connections, then there are two ways to handle this.

The first option is to explicitly choose a connection for each SQL statement, for example:

```
EXEC SQL AT connection-name SELECT ...;
```

This option is particularly suitable if the application needs to use several connections in mixed order. If your application uses multiple threads of execution, they cannot share a connection concurrently. You must either explicitly control access to the connection (using mutexes) or use a connection for each thread.

The second option is to execute a statement to switch the current connection. That statement is:

```
EXEC SQL SET CONNECTION connection-name;
```

This option is particularly convenient if many statements are to be executed on the same connection. Here is an example program managing multiple database connections:

```c
#include <stdio.h>

EXEC SQL BEGIN DECLARE SECTION;
  char dbname[1024];
EXEC SQL END DECLARE SECTION;

int main()
{
  EXEC SQL CONNECT TO testdb1 AS con1 USER testuser;
  EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
  EXEC SQL CONNECT TO testdb2 AS con2 USER testuser;
  EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
  EXEC SQL CONNECT TO testdb3 AS con3 USER testuser;
  EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

  /* This query would be executed in the last opened database "testdb3". */
  EXEC SQL SELECT current_database() INTO :dbname;
  printf("current=%s (should be testdb3)\n", dbname);

  /* Using "AT" to run a query in "testdb2" */
  EXEC SQL AT con2 SELECT current_database() INTO :dbname;
  printf("current=%s (should be testdb2)\n", dbname);

  /* Switch the current connection to "testdb1". */
  EXEC SQL SET CONNECTION con1;
}```
EXEC SQL SELECT current_database() INTO :dbname;
printf("current=%s (should be testdb1)\n", dbname);

EXEC SQL DISCONNECT ALL;
return 0;
}

This example would produce this output:
current=testdb3 (should be testdb3)
current=testdb2 (should be testdb2)
current=testdb1 (should be testdb1)

34.2.3. Closing a Connection

To close a connection, use the following statement:

EXEC SQL DISCONNECT [connection];

The connection can be specified in the following ways:

- connection-name
- DEFAULT
- CURRENT
- ALL

If no connection name is specified, the current connection is closed.

It is good style that an application always explicitly disconnect from every connection it opened.

34.3. Running SQL Commands

Any SQL command can be run from within an embedded SQL application. Below are some examples of how to do that.

34.3.1. Executing SQL Statements

Creating a table:

EXEC SQL CREATE TABLE foo (number integer, ascii char(16));
EXEC SQL CREATE UNIQUE INDEX num1 ON foo(number);
EXEC SQL COMMIT;

Inserting rows:

EXEC SQL INSERT INTO foo (number, ascii) VALUES (9999, 'doodad');
EXEC SQL COMMIT;

Deleting rows:
EXEC SQL DELETE FROM foo WHERE number = 9999;
EXEC SQL COMMIT;

Updates:
EXEC SQL UPDATE foo
    SET ascii = 'foobar'
    WHERE number = 9999;
EXEC SQL COMMIT;

SELECT statements that return a single result row can also be executed using EXEC SQL directly. To handle result sets with multiple rows, an application has to use a cursor; see Section 34.3.2 below. (As a special case, an application can fetch multiple rows at once into an array host variable; see Section 34.4.4.3.1.)

Single-row select:
EXEC SQL SELECT foo INTO :FooBar FROM table1 WHERE ascii = 'doodad';

Also, a configuration parameter can be retrieved with the SHOW command:
EXEC SQL SHOW search_path INTO :var;

The tokens of the form :something are host variables, that is, they refer to variables in the C program. They are explained in Section 34.4.

34.3.2. Using Cursors

To retrieve a result set holding multiple rows, an application has to declare a cursor and fetch each row from the cursor. The steps to use a cursor are the following: declare a cursor, open it, fetch a row from the cursor, repeat, and finally close it.

Select using cursors:
EXEC SQL DECLARE foo_bar CURSOR FOR
    SELECT number, ascii FROM foo
    ORDER BY ascii;
EXEC SQL OPEN foo_bar;
EXEC SQL FETCH foo_bar INTO :FooBar, DooDad;
...
EXEC SQL CLOSE foo_bar;
EXEC SQL COMMIT;
For more details about declaration of the cursor, see DECLARE, and see FETCH for FETCH command details.

Note: The ECPG DECLARE command does not actually cause a statement to be sent to the PostgreSQL backend. The cursor is opened in the backend (using the backend’s DECLARE command) at the point when the OPEN command is executed.

34.3.3. Managing Transactions

In the default mode, statements are committed only when EXEC SQL COMMIT is issued. The embedded SQL interface also supports autocommit of transactions (similar to psql’s default behavior) via the -t command-line option to ecpg (see ecpg) or via the EXEC SQL SET AUTOCOMMIT TO ON statement. In autocommit mode, each command is automatically committed unless it is inside an explicit transaction block. This mode can be explicitly turned off using EXEC SQL SET AUTOCOMMIT TO OFF.

The following transaction management commands are available:

EXEC SQL COMMIT

Commit an in-progress transaction.

EXEC SQL ROLLBACK

Roll back an in-progress transaction.

EXEC SQL SET AUTOCOMMIT TO ON

Enable autocommit mode.

SET AUTOCOMMIT TO OFF

Disable autocommit mode. This is the default.

34.3.4. Prepared Statements

When the values to be passed to an SQL statement are not known at compile time, or the same statement is going to be used many times, then prepared statements can be useful.

The statement is prepared using the command PREPARE. For the values that are not known yet, use the placeholder “?”:

EXEC SQL PREPARE stmt1 FROM "SELECT oid, datname FROM pg_database WHERE oid = ?";

If a statement returns a single row, the application can call EXECUTE after PREPARE to execute the statement, supplying the actual values for the placeholders with a USING clause:

EXEC SQL EXECUTE stmt1 INTO :dboid, :dbname USING 1;

If a statement returns multiple rows, the application can use a cursor declared based on the prepared statement. To bind input parameters, the cursor must be opened with a USING clause:
EXEC SQL PREPARE stmt1 FROM "SELECT oid, datname FROM pg_database WHERE oid > ?";
EXEC SQL DECLARE foo_bar CURSOR FOR stmt1;

/* when end of result set reached, break out of while loop */
EXEC SQL WHENEVER NOT FOUND DO BREAK;

EXEC SQL OPEN foo_bar USING 100;
...
while (1) {
    EXEC SQL FETCH NEXT FROM foo_bar INTO :dboid, :dbname;
    ...
}
EXEC SQL CLOSE foo_bar;

When you don’t need the prepared statement anymore, you should deallocate it:

EXEC SQL DEALLOCATE PREPARE name;

For more details about PREPARE, see PREPARE. Also see Section 34.5 for more details about using placeholders and input parameters.

### 34.4. Using Host Variables

In Section 34.3 you saw how you can execute SQL statements from an embedded SQL program. Some of those statements only used fixed values and did not provide a way to insert user-supplied values into statements or have the program process the values returned by the query. Those kinds of statements are not really useful in real applications. This section explains in detail how you can pass data between your C program and the embedded SQL statements using a simple mechanism called host variables. In an embedded SQL program we consider the SQL statements to be guests in the C program code which is the host language. Therefore the variables of the C program are called host variables.

Another way to exchange values between PostgreSQL backends and ECPG applications is the use of SQL descriptors, described in Section 34.7.

#### 34.4.1. Overview

Passing data between the C program and the SQL statements is particularly simple in embedded SQL. Instead of having the program paste the data into the statement, which entails various complications, such as properly quoting the value, you can simply write the name of a C variable into the SQL statement, prefixed by a colon. For example:

EXEC SQL INSERT INTO sometable VALUES (:v1, ‘foo’, :v2);

This statement refers to two C variables named v1 and v2 and also uses a regular SQL string literal, to illustrate that you are not restricted to use one kind of data or the other.

This style of inserting C variables in SQL statements works anywhere a value expression is expected in an SQL statement.
34.4.2. Declare Sections

To pass data from the program to the database, for example as parameters in a query, or to pass data from the database back to the program, the C variables that are intended to contain this data need to be declared in specially marked sections, so the embedded SQL preprocessor is made aware of them. This section starts with:

EXEC SQL BEGIN DECLARE SECTION;

and ends with:

EXEC SQL END DECLARE SECTION;

Between those lines, there must be normal C variable declarations, such as:

```c
int x = 4;
char foo[16], bar[16];
```

As you can see, you can optionally assign an initial value to the variable. The variable’s scope is determined by the location of its declaring section within the program. You can also declare variables with the following syntax which implicitly creates a declare section:

```c
EXEC SQL int i = 4;
```

You can have as many declare sections in a program as you like.

The declarations are also echoed to the output file as normal C variables, so there’s no need to declare them again. Variables that are not intended to be used in SQL commands can be declared normally outside these special sections.

The definition of a structure or union also must be listed inside a DECLARE section. Otherwise the preprocessor cannot handle these types since it does not know the definition.

34.4.3. Retrieving Query Results

Now you should be able to pass data generated by your program into an SQL command. But how do you retrieve the results of a query? For that purpose, embedded SQL provides special variants of the usual commands SELECT and FETCH. These commands have a special INTO clause that specifies which host variables the retrieved values are to be stored in. SELECT is used for a query that returns only single row, and FETCH is used for a query that returns multiple rows, using a cursor.

Here is an example:

```c
/*
 * assume this table:
 * CREATE TABLE test1 (a int, b varchar(50));
 */

EXEC SQL BEGIN DECLARE SECTION;
int v1;
VARCHAR v2;
EXEC SQL END DECLARE SECTION;

...
Chapter 34. ECPG - Embedded SQL in C

So the INTO clause appears between the select list and the FROM clause. The number of elements in the select list and the list after INTO (also called the target list) must be equal.

Here is an example using the command FETCH:

```sql
EXEC SQL BEGIN DECLARE SECTION;
  int v1;
  VARCHAR v2;
EXEC SQL END DECLARE SECTION;
...
EXEC SQL DECLARE foo CURSOR FOR SELECT a, b FROM test;
...
do {
  ...
    EXEC SQL FETCH NEXT FROM foo INTO :v1, :v2;
    ...
} while (...);
```

Here the INTO clause appears after all the normal clauses.

34.4.4. Type Mapping

When ECPG applications exchange values between the PostgreSQL server and the C application, such as when retrieving query results from the server or executing SQL statements with input parameters, the values need to be converted between PostgreSQL data types and host language variable types (C language data types, concretely). One of the main points of ECPG is that it takes care of this automatically in most cases.

In this respect, there are two kinds of data types: Some simple PostgreSQL data types, such as integer and text, can be read and written by the application directly. Other PostgreSQL data types, such as timestamp and numeric can only be accessed through special library functions; see Section 34.4.4.2.

Table 34-1 shows which PostgreSQL data types correspond to which C data types. When you wish to send or receive a value of a given PostgreSQL data type, you should declare a C variable of the corresponding C data type in the declare section.

Table 34-1. Mapping Between PostgreSQL Data Types and C Variable Types

<table>
<thead>
<tr>
<th>PostgreSQL data type</th>
<th>Host variable type</th>
</tr>
</thead>
<tbody>
<tr>
<td>smallint</td>
<td>short</td>
</tr>
<tr>
<td>integer</td>
<td>int</td>
</tr>
<tr>
<td>bigint</td>
<td>long long int</td>
</tr>
<tr>
<td>decimal</td>
<td>decimal</td>
</tr>
<tr>
<td>numeric</td>
<td>numeric</td>
</tr>
<tr>
<td>real</td>
<td>float</td>
</tr>
<tr>
<td>double precision</td>
<td>double</td>
</tr>
<tr>
<td>smallserial</td>
<td>short</td>
</tr>
</tbody>
</table>
### 34.4.4.1. Handling Character Strings

To handle SQL character string data types, such as `varchar` and `text`, there are two possible ways to declare the host variables.

One way is using `char[]`, an array of `char`, which is the most common way to handle character data in C.

```sql
EXEC SQL BEGIN DECLARE SECTION;
    char str[50];
EXEC SQL END DECLARE SECTION;
```

Note that you have to take care of the length yourself. If you use this host variable as the target variable of a query which returns a string with more than 49 characters, a buffer overflow occurs.

The other way is using the `VARCHAR` type, which is a special type provided by ECPG. The definition on an array of type `VARCHAR` is converted into a named `struct` for every variable. A declaration like:

```sql
VARCHAR var[180];
```

is converted into:

```c
struct varchar_var { int len; char arr[180]; } var;
```

The member `arr` hosts the string including a terminating zero byte. Thus, to store a string in a `VARCHAR` host variable, the host variable has to be declared with the length including the zero byte terminator. The member `len` holds the length of the string stored in the `arr` without the terminating zero byte. When a host variable is used as input for a query, if `strlen(arr)` and `len` are different, the shorter one is used.

`VARCHAR` can be written in upper or lower case, but not in mixed case.

`char` and `VARCHAR` host variables can also hold values of other SQL types, which will be stored in their string forms.
34.4.4.2. Accessing Special Data Types

ECPG contains some special types that help you to interact easily with some special data types from the PostgreSQL server. In particular, it has implemented support for the numeric, decimal, date, timestamp, and interval types. These data types cannot usefully be mapped to primitive host variable types (such as int, long long int, or char[]), because they have a complex internal structure. Applications deal with these types by declaring host variables in special types and accessing them using functions in the pgtypes library. The pgtypes library, described in detail in Section 34.6 contains basic functions to deal with those types, such that you do not need to send a query to the SQL server just for adding an interval to a time stamp for example.

The follow subsections describe these special data types. For more details about pgtypes library functions, see Section 34.6.

34.4.4.2.1. timestamp, date

Here is a pattern for handling timestamp variables in the ECPG host application.

First, the program has to include the header file for the timestamp type:

```
#include <pgtypes_timestamp.h>
```

Next, declare a host variable as type timestamp in the declare section:

```
EXEC SQL BEGIN DECLARE SECTION;
timestamp ts;
EXEC SQL END DECLARE SECTION;
```

And after reading a value into the host variable, process it using pgtypes library functions. In following example, the timestamp value is converted into text (ASCII) form with the PGTYPEStimestamp_to_asc() function:

```
EXEC SQL SELECT now()::timestamp INTO :ts;
printf("ts = %s\n", PGTYPEStimestamp_to_asc(ts));
```

This example will show some result like following:

```
ts = 2010-06-27 18:03:56.949343
```

In addition, the DATE type can be handled in the same way. The program has to include pgtypes_date.h, declare a host variable as the date type and convert a DATE value into a text form using PGTYPESdate_to_asc() function. For more details about the pgtypes library functions, see Section 34.6.

34.4.4.2.2. interval

The handling of the interval type is also similar to the timestamp and date types. It is required, however, to allocate memory for an interval type value explicitly. In other words, the memory space for the variable has to be allocated in the heap memory, not in the stack memory.

Here is an example program:
#include <stdio.h>
#include <stdlib.h>
#include <pgtypes_interval.h>

int main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    interval *in;
    EXEC SQL END DECLARE SECTION;

    EXEC SQL CONNECT TO testdb;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

    in = PGTYPESinterval_new();
    EXEC SQL SELECT '1 min'::interval INTO :in;
    printf("interval = %s\n", PGTYPESinterval_to_asc(in));
    PGTYPESinterval_free(in);

    EXEC SQL COMMIT;
    EXEC SQL DISCONNECT ALL;
    return 0;
}

34.4.4.2.3. numeric, decimal

The handling of the numeric and decimal types is similar to the interval type: It requires defining a pointer, allocating some memory space on the heap, and accessing the variable using the pgtypes library functions. For more details about the pgtypes library functions, see Section 34.6.

No functions are provided specifically for the decimal type. An application has to convert it to a numeric variable using a pgtypes library function to do further processing.

Here is an example program handling numeric and decimal type variables.

#include <stdio.h>
#include <stdlib.h>
#include <pgtypes_numeric.h>

EXEC SQL WHENEVER SQLERROR STOP;

int main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    numeric *num;
    numeric *num2;
    decimal *dec;
    EXEC SQL END DECLARE SECTION;

    EXEC SQL CONNECT TO testdb;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

    num = PGTYPESNumeric_new();
    dec = PGTYPESDecimal_new();
EXEC SQL SELECT 12.345::numeric(4,2), 23.456::decimal(4,2) INTO :num, :dec;

printf("numeric = %s\n", PGTYPESnumeric_to_asc(num, 0));
printf("numeric = %s\n", PGTYPESnumeric_to_asc(num, 1));
printf("numeric = %s\n", PGTYPESnumeric_to_asc(num, 2));

/* Convert decimal to numeric to show a decimal value. */
num2 = PGTYPESnumeric_new();
PGTYPESnumeric_from_decimal(dec, num2);

printf("decimal = %s\n", PGTYPESnumeric_to_asc(num2, 0));
printf("decimal = %s\n", PGTYPESnumeric_to_asc(num2, 1));
printf("decimal = %s\n", PGTYPESnumeric_to_asc(num2, 2));

PGTYPESnumeric_free(num2);
PGTYPESdecimal_free(dec);
PGTYPESnumeric_free(num);

EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;
return 0;
}

34.4.4.3. Host Variables with Nonprimitive Types

As a host variable you can also use arrays, typedefs, structs, and pointers.

34.4.4.3.1. Arrays

There are two use cases for arrays as host variables. The first is a way to store some text string in char[] or VARCHAR[], as explained in Section 34.4.4.1. The second use case is to retrieve multiple rows from a query result without using a cursor. Without an array, to process a query result consisting of multiple rows, it is required to use a cursor and the FETCH command. But with array host variables, multiple rows can be received at once. The length of the array has to be defined to be able to accommodate all rows, otherwise a buffer overflow will likely occur.

Following example scans the pg_database system table and shows all OIDs and names of the available databases:

int
main(void)
{
EXEC SQL BEGIN DECLARE SECTION;
  int dbid[8];
  char dbname[8][16];
  int i;
EXEC SQL END DECLARE SECTION;

memset(dbname, 0, sizeof(char) * 16 * 8);
memset(dbid, 0, sizeof(int) * 8);

EXEC SQL CONNECT TO testdb;
Chapter 34. ECPG - Embedded SQL in C

EXEC SQL SElECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;

/* Retrieve multiple rows into arrays at once. */
EXEC SQL SELECT oid, datname INTO :dbid, :dbname FROM pg_database;

for (i = 0; i < 8; i++)
    printf("oid=%d, dbname=%s\n", dbid[i], dbname[i]);

EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;
return 0;
}

This example shows following result. (The exact values depend on local circumstances.)

oid=1, dbname=template1
oid=11510, dbname=template0
oid=11511, dbname=postgres
oid=313780, dbname=testdb
oid=0, dbname=
oid=0, dbname=
oid=0, dbname=

34.4.4.3.2. Structures

A structure whose member names match the column names of a query result, can be used to retrieve multiple columns at once. The structure enables handling multiple column values in a single host variable.

The following example retrieves OIDs, names, and sizes of the available databases from the pg_database system table and using the pg_database_size() function. In this example, a structure variable dbinfo_t with members whose names match each column in the SELECT result is used to retrieve one result row without putting multiple host variables in the FETCH statement.

EXEC SQL BEGIN DECLARE SECTION;
typedef struct
{
    int oid;
    char datname[65];
    long long int size;
} dbinfo_t;

dbinfo_t dbval;
EXEC SQL END DECLARE SECTION;

memset(&dbval, 0, sizeof(dbinfo_t));

EXEC SQL DECLARE cur1 CURSOR FOR SELECT oid, datname, pg_database_size(oid) AS size FROM pg_database;
EXEC SQL OPEN cur1;

/* when end of result set reached, break out of while loop */
EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1)
    
826
/* Fetch multiple columns into one structure. */
EXEC SQL FETCH FROM cur1 INTO :dbval;

/* Print members of the structure. */
printf("oid=%d, datname=%s, size=%lld\n", dbval.oid, dbval.datname, dbval.size);
}

EXEC SQL CLOSE cur1;

This example shows following result. (The exact values depend on local circumstances.)

oid=1, datname=template1, size=4324580
oid=11510, datname=template0, size=4243460
oid=11511, datname=postgres, size=4324580
oid=313780, datname=testdb, size=8183012

Structure host variables “absorb” as many columns as the structure as fields. Additional columns can be assigned to other host variables. For example, the above program could also be restructured like this, with the size variable outside the structure:

EXEC SQL BEGIN DECLARE SECTION;

typedef struct
{
    int oid;
    char datname[65];
} dbinfo_t;

dbinfo_t dbval;
    long long int size;
EXEC SQL END DECLARE SECTION;

memset(&dbval, 0, sizeof(dbinfo_t));

EXEC SQL DECLARE cur1 CURSOR FOR SELECT oid, datname, pg_database_size(oid) AS size FROM pg_database;
EXEC SQL OPEN cur1;

/* when end of result set reached, break out of while loop */
EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1)
{
    /* Fetch multiple columns into one structure. */
    EXEC SQL FETCH FROM cur1 INTO :dbval, :size;

    /* Print members of the structure. */
    printf("oid=%d, datname=%s, size=%lld\n", dbval.oid, dbval.datname, size);
}

EXEC SQL CLOSE cur1;

827
34.4.3.3. Typedefs

Use the `typedef` keyword to map new types to already existing types.

```sql
EXEC SQL BEGIN DECLARE SECTION;
  typedef char mychartype[40];
  typedef long serial_t;
EXEC SQL END DECLARE SECTION;
```

Note that you could also use:

```sql
EXEC SQL TYPE serial_t IS long;
```

This declaration does not need to be part of a declare section.

34.4.4.3.4. Pointers

You can declare pointers to the most common types. Note however that you cannot use pointers as target variables of queries without auto-allocation. See Section 34.7 for more information on auto-allocation.

```sql
EXEC SQL BEGIN DECLARE SECTION;
  int *intp;
  char **charp;
EXEC SQL END DECLARE SECTION;
```

34.4.5. Handling Nonprimitive SQL Data Types

This section contains information on how to handle nonscalar and user-defined SQL-level data types in ECPG applications. Note that this is distinct from the handling of host variables of nonprimitive types, described in the previous section.

34.4.5.1. Arrays

Multi-dimensional SQL-level arrays are not directly supported in ECPG. One-dimensional SQL-level arrays can be mapped into C array host variables and vice-versa. However, when creating a statement `ecpg` does not know the types of the columns, so that it cannot check if a C array is input into a corresponding SQL-level array. When processing the output of a SQL statement, `ecpg` has the necessary information and thus checks if both are arrays.

If a query accesses elements of an array separately, then this avoids the use of arrays in ECPG. Then, a host variable with a type that can be mapped to the element type should be used. For example, if a column type is array of integer, a host variable of type `int` can be used. Also if the element type is `varchar` or `text`, a host variable of type `char[]` or `VARCHAR[]` can be used.

Here is an example. Assume the following table:

```sql
CREATE TABLE t3 (  
  ii integer[]
);```
testdb=> SELECT * FROM t3;
   ii
  (1,2,3,4,5)
  (1 row)

The following example program retrieves the 4th element of the array and stores it into a host variable of type `int`:

```sql
EXEC SQL BEGIN DECLARE SECTION;
int ii;
EXEC SQL END DECLARE SECTION;
EXEC SQL DECLARE cur1 CURSOR FOR SELECT ii[4] FROM t3;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1)
{
   EXEC SQL FETCH FROM cur1 INTO :ii ;
   printf("ii=%d\n", ii);
}
EXEC SQL CLOSE cur1;
```

This example shows the following result:

```
ii=4
```

To map multiple array elements to the multiple elements in an array type host variables each element of array column and each element of the host variable array have to be managed separately, for example:

```sql
EXEC SQL BEGIN DECLARE SECTION;
int ii_a[8];
EXEC SQL END DECLARE SECTION;
EXEC SQL DECLARE cur1 CURSOR FOR SELECT ii[1], ii[2], ii[3], ii[4] FROM t3;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1)
{
   EXEC SQL FETCH FROM cur1 INTO :ii_a[0], :ii_a[1], :ii_a[2], :ii_a[3];
   ...
}
```

Note again that

```sql
EXEC SQL BEGIN DECLARE SECTION;
int ii_a[8];
EXEC SQL END DECLARE SECTION;
```
EXEC SQL DECLARE cur1 CURSOR FOR SELECT ii FROM t3;
EXEC SQL OPEN cur1;

EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1) {
    /* WRONG */
    EXEC SQL FETCH FROM cur1 INTO :ii_a;
    ...
}

would not work correctly in this case, because you cannot map an array type column to an array host variable directly.

Another workaround is to store arrays in their external string representation in host variables of type char[] or VARCHAR[]. For more details about this representation, see Section 8.15.2. Note that this means that the array cannot be accessed naturally as an array in the host program (without further processing that parses the text representation).

34.4.5.2. Composite Types

Composite types are not directly supported in ECPG, but an easy workaround is possible. The available workarounds are similar to the ones described for arrays above: Either access each attribute separately or use the external string representation.

For the following examples, assume the following type and table:

CREATE TYPE comp_t AS (intval integer, textval varchar(32));
CREATE TABLE t4 (compval comp_t);
INSERT INTO t4 VALUES ( (256, 'PostgreSQL') );

The most obvious solution is to access each attribute separately. The following program retrieves data from the example table by selecting each attribute of the type comp_t separately:

EXEC SQL BEGIN DECLARE SECTION;
int intval;
varchar textval[33];
EXEC SQL END DECLARE SECTION;

/* Put each element of the composite type column in the SELECT list. */
EXEC SQL DECLARE cur1 CURSOR FOR SELECT (compval).intval, (compval).textval FROM t4;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1) {
    /* Fetch each element of the composite type column into host variables. */
    EXEC SQL FETCH FROM cur1 INTO :intval, :textval;
    printf("intval=%d, textval=%s\n", intval, textval.arr);
}

EXEC SQL CLOSE cur1;
To enhance this example, the host variables to store values in the `FETCH` command can be gathered into one structure. For more details about the host variable in the structure form, see Section 34.4.4.3.2. To switch to the structure, the example can be modified as below. The two host variables, `intval` and `textval`, become members of the `comp_t` structure, and the structure is specified on the `FETCH` command.

```c
EXEC SQL BEGIN DECLARE SECTION;
typedef struct {
    int intval;
    varchar textval[33];
} comp_t;

comp_t compval;
EXEC SQL END DECLARE SECTION;

EXEC SQL DECLARE cur1 CURSOR FOR SELECT (compval).intval, (compval).textval FROM t4;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1) {
    /* Put all values in the SELECT list into one structure. */
    EXEC SQL FETCH FROM cur1 INTO :compval;
    printf("intval=%d, textval=%s\n", compval.intval, compval.textval.arr);
}
EXEC SQL CLOSE cur1;
```

Although a structure is used in the `FETCH` command, the attribute names in the `SELECT` clause are specified one by one. This can be enhanced by using a `*` to ask for all attributes of the composite type value.

```c
EXEC SQL DECLARE cur1 CURSOR FOR SELECT (compval).* FROM t4;
EXEC SQL OPEN cur1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1) {
    /* Put all values in the SELECT list into one structure. */
    EXEC SQL FETCH FROM cur1 INTO :compval;
    printf("intval=%d, textval=%s\n", compval.intval, compval.textval.arr);
}
```

This way, composite types can be mapped into structures almost seamlessly, even though ECPG does not understand the composite type itself.

Finally, it is also possible to store composite type values in their external string representation in host variables of type `char[]` or `VARCHAR[]`. But that way, it is not easily possible to access the fields of the value from the host program.
Chapter 34. ECPG - Embedded SQL in C

34.4.5.3. User-defined Base Types

New user-defined base types are not directly supported by ECPG. You can use the external string representation and host variables of type char[] or VARCHAR[], and this solution is indeed appropriate and sufficient for many types.

Here is an example using the data type complex from the example in Section 36.11. The external string representation of that type is (%f,%f), which is defined in the functions complex_in() and complex_out() functions in Section 36.11. The following example inserts the complex type values (1,1) and (3,3) into the columns a and b, and select them from the table after that.

EXEC SQL BEGIN DECLARE SECTION;
    varchar a[64];
    varchar b[64];
EXEC SQL END DECLARE SECTION;

EXEC SQL INSERT INTO test_complex VALUES ('(1,1)', '(3,3)');

EXEC SQL DECLARE cur1 CURSOR FOR SELECT a, b FROM test_complex;
EXEC SQL OPEN cur1;

EXEC SQL WHENEVER NOT FOUND DO BREAK;

while (1) {
    EXEC SQL FETCH FROM cur1 INTO :a, :b;
    printf("a=%s, b=%s\n", a.arr, b.arr);
}

EXEC SQL CLOSE cur1;

This example shows following result:

a=(1,1), b=(3,3)

Another workaround is avoiding the direct use of the user-defined types in ECPG and instead create a function or cast that converts between the user-defined type and a primitive type that ECPG can handle. Note, however, that type casts, especially implicit ones, should be introduced into the type system very carefully.

For example,

CREATE FUNCTION create_complex(r double, i double) RETURNS complex
    LANGUAGE SQL
    IMMUTABLE
AS $$ SELECT $1 * complex '(1,0)' + $2 * complex '(0,1)' $$;

After this definition, the following

EXEC SQL BEGIN DECLARE SECTION;
    double a, b, c, d;
EXEC SQL END DECLARE SECTION;

    a = 1;
    b = 2;
    c = 3;
    d = 4;
Chapter 34. ECPG - Embedded SQL in C

EXEC SQL INSERT INTO test_complex VALUES (create_complex(:a, :b), create_complex(:c, :d))

has the same effect as

EXEC SQL INSERT INTO test_complex VALUES ('(1,2)', '(3,4)');

34.4.6. Indicators

The examples above do not handle null values. In fact, the retrieval examples will raise an error if they fetch a null value from the database. To be able to pass null values to the database or retrieve null values from the database, you need to append a second host variable specification to each host variable that contains data. This second host variable is called the indicator and contains a flag that tells whether the datum is null, in which case the value of the real host variable is ignored. Here is an example that handles the retrieval of null values correctly:

EXEC SQL BEGIN DECLARE SECTION;
VARCHAR val;
int val_ind;
EXEC SQL END DECLARE SECTION:

... EXEC SQL SELECT b INTO :val :val_ind FROM test1;

The indicator variable val_ind will be zero if the value was not null, and it will be negative if the value was null.

The indicator has another function: if the indicator value is positive, it means that the value is not null, but it was truncated when it was stored in the host variable.

If the argument -r no_indicator is passed to the preprocessor ecpg, it works in "no-indicator" mode. In no-indicator mode, if no indicator variable is specified, null values are signaled (on input and output) for character string types as empty string and for integer types as the lowest possible value for type (for example, INT_MIN for int).

34.5. Dynamic SQL

In many cases, the particular SQL statements that an application has to execute are known at the time the application is written. In some cases, however, the SQL statements are composed at run time or provided by an external source. In these cases you cannot embed the SQL statements directly into the C source code, but there is a facility that allows you to call arbitrary SQL statements that you provide in a string variable.
34.5.1. Executing Statements without a Result Set

The simplest way to execute an arbitrary SQL statement is to use the command `EXECUTE IMMEDIATE`. For example:

```c
EXEC SQL BEGIN DECLARE SECTION;
const char *stmt = "CREATE TABLE test1 (...);
EXEC SQL END DECLARE SECTION;
EXEC SQL EXECUTE IMMEDIATE :stmt;
```

`EXECUTE IMMEDIATE` can be used for SQL statements that do not return a result set (e.g., DDL, `INSERT`, `UPDATE`, `DELETE`). You cannot execute statements that retrieve data (e.g., `SELECT`) this way. The next section describes how to do that.

34.5.2. Executing a Statement with Input Parameters

A more powerful way to execute arbitrary SQL statements is to prepare them once and execute the prepared statement as often as you like. It is also possible to prepare a generalized version of a statement and then execute specific versions of it by substituting parameters. When preparing the statement, write question marks where you want to substitute parameters later. For example:

```c
EXEC SQL BEGIN DECLARE SECTION;
const char *stmt = "INSERT INTO test1 VALUES(?, ?);
EXEC SQL END DECLARE SECTION;
EXEC SQL PREPARE mystmt FROM :stmt;
...
EXEC SQL EXECUTE mystmt USING 42, 'foobar';
```

When you don’t need the prepared statement anymore, you should deallocate it:

```c
EXEC SQL DEALLOCATE PREPARE name;
```

34.5.3. Executing a Statement with a Result Set

To execute an SQL statement with a single result row, `EXECUTE` can be used. To save the result, add an `INTO` clause.

```c
EXEC SQL BEGIN DECLARE SECTION;
const char *stmt = "SELECT a, b, c FROM test1 WHERE a > ?";
int v1, v2;
VARCHAR v3[50];
EXEC SQL END DECLARE SECTION;
EXEC SQL PREPARE mystmt FROM :stmt;
...
EXEC SQL EXECUTE mystmt INTO :v1, :v2, :v3 USING 37;
```

An `EXECUTE` command can have an `INTO` clause, a `USING` clause, both, or neither.
If a query is expected to return more than one result row, a cursor should be used, as in the following example. (See Section 34.3.2 for more details about the cursor.)

```sql
EXEC SQL BEGIN DECLARE SECTION;
    char dbaname[128];
    char datname[128];
    char *stmt = "SELECT u.usename as dbaname, d.datname "
                 " FROM pg_database d, pg_user u "
                 " WHERE d.datdba = u.usesysid";
EXEC SQL END DECLARE SECTION;
EXEC SQL CONNECT TO testdb AS con1 USER testuser;
EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;
EXEC SQL PREPARE stmt1 FROM :stmt;
EXEC SQL DECLARE cursor1 CURSOR FOR stmt1;
EXEC SQL OPEN cursor1;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
while (1) {
    EXEC SQL FETCH cursor1 INTO :dbaname, :datname;
    printf("dbaname=%s, datname=%s\n", dbaname, datname);
}
EXEC SQL CLOSE cursor1;
EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;
```

### 34.6. pgtypes Library

The pgtypes library maps PostgreSQL database types to C equivalents that can be used in C programs. It also offers functions to do basic calculations with those types within C, i.e., without the help of the PostgreSQL server. See the following example:

```sql
EXEC SQL BEGIN DECLARE SECTION;
    date date1;
    timestamp ts1, tsout;
    interval iv1;
    char *out;
EXEC SQL END DECLARE SECTION;
PGTYPESdate_today(&date1);
EXEC SQL SELECT started, duration INTO :ts1, :iv1 FROM datetbl WHERE d=:date1;
PGTYPEStimestamp_add_interval(&ts1, &iv1, &tsout);
out = PGTYPEStimestamp_to_asc(&tsout);
printf("Started + duration: %s\n", out);
PGTYPESchar_free(out);
```
34.6.1. Character Strings

Some functions such as `PGTYPEnumeric_to_asc` return a pointer to a freshly allocated character string. These results should be freed with `PGTYPEchar_free` instead of `free`. (This is important only on Windows, where memory allocation and release sometimes need to be done by the same library.)

34.6.2. The numeric Type

The numeric type offers to do calculations with arbitrary precision. See Section 8.1 for the equivalent type in the PostgreSQL server. Because of the arbitrary precision this variable needs to be able to expand and shrink dynamically. That’s why you can only create numeric variables on the heap, by means of the `PGTYPEnumeric_new` and `PGTYPEnumeric_free` functions. The decimal type, which is similar but limited in precision, can be created on the stack as well as on the heap.

The following functions can be used to work with the numeric type:

`PGTYPEnumeric_new`

Request a pointer to a newly allocated numeric variable.
```c
numeric *PGTYPEnumeric_new(void);
```

`PGTYPEnumeric_free`

Free a numeric type, release all of its memory.
```c
void PGTYPESnumeric_free(numeric *var);
```

`PGTYPEnumeric_from_asc`

Parse a numeric type from its string notation.
```c
numeric *PGTYPEnumeric_from_asc(char *str, char **endptr);
```

Valid formats are for example: `-2.794, .794, +3.44, 592.49E07` or `-32.84e-4`. If the value could be parsed successfully, a valid pointer is returned, else the NULL pointer. At the moment ECPG always parses the complete string and so it currently does not support to store the address of the first invalid character in `*endptr`. You can safely set `endptr` to NULL.

`PGTYPEnumeric_to_asc`

Returns a pointer to a string allocated by `malloc` that contains the string representation of the numeric type `num`.
```c
char *PGTYPEnumeric_to_asc(numeric *num, int dscale);
```

The numeric value will be printed with `dscale` decimal digits, with rounding applied if necessary. The result must be freed with `PGTYPEchar_free()`.

`PGTYPEnumeric_add`

Add two numeric variables into a third one.
```c
int PGTYPESnumeric_add(numeric *var1, numeric *var2, numeric *result);
```

The function adds the variables `var1` and `var2` into the result variable `result`. The function returns 0 on success and -1 in case of error.
Chapter 34. ECPG - Embedded SQL in C

PGTYPEnumeric_sub

Subtract two numeric variables and return the result in a third one.

int PGTYPESnumeric_sub(numeric *var1, numeric *var2, numeric *result);
The function subtracts the variable \( \text{var2} \) from the variable \( \text{var1} \). The result of the operation is stored in the variable \( \text{result} \). The function returns 0 on success and -1 in case of error.

PGTYPEnumeric_mul

Multiply two numeric variables and return the result in a third one.

int PGTYPESnumeric_mul(numeric *var1, numeric *var2, numeric *result);
The function multiplies the variables \( \text{var1} \) and \( \text{var2} \). The result of the operation is stored in the variable \( \text{result} \). The function returns 0 on success and -1 in case of error.

PGTYPEnumeric_div

Divide two numeric variables and return the result in a third one.

int PGTYPESnumeric_div(numeric *var1, numeric *var2, numeric *result);
The function divides the variables \( \text{var1} \) by \( \text{var2} \). The result of the operation is stored in the variable \( \text{result} \). The function returns 0 on success and -1 in case of error.

PGTYPEnumeric_cmp

Compare two numeric variables.

int PGTYPESnumeric_cmp(numeric *var1, numeric *var2);
This function compares two numeric variables. In case of error, \( \text{INT\_MAX} \) is returned. On success, the function returns one of three possible results:

- 1, if \( \text{var1} \) is bigger than \( \text{var2} \)
- -1, if \( \text{var1} \) is smaller than \( \text{var2} \)
- 0, if \( \text{var1} \) and \( \text{var2} \) are equal

PGTYPEnumeric_from_int

Convert an int variable to a numeric variable.

int PGTYPESnumeric_from_int(signed int int_val, numeric *var);
This function accepts a variable of type signed int and stores it in the numeric variable \( \text{var} \). Upon success, 0 is returned and -1 in case of a failure.

PGTYPEnumeric_from_long

Convert a long int variable to a numeric variable.

int PGTYPESnumeric_from_long(signed long int long_val, numeric *var);
This function accepts a variable of type signed long int and stores it in the numeric variable \( \text{var} \). Upon success, 0 is returned and -1 in case of a failure.

PGTYPEnumeric_copy

Copy over one numeric variable into another one.

int PGTYPESnumeric_copy(numeric *src, numeric *dst);
This function copies over the value of the variable that \( \text{src} \) points to into the variable that \( \text{dst} \) points to. It returns 0 on success and -1 if an error occurs.

PGTYPEnumeric_from_double

Convert a variable of type double to a numeric.

int PGTYPESnumeric_from_double(double d, numeric *dst);
This function accepts a variable of type double and stores the result in the variable that dst points to. It returns 0 on success and -1 if an error occurs.

**PGTYPESnumeric_to_double**

Convert a variable of type numeric to double.

```c
int PGTYPESnumeric_to_double(numeric *nv, double *dp)
```

The function converts the numeric value from the variable that nv points to into the double variable that dp points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable errno will be set to PGTYPES_NUM_OVERFLOW additionally.

**PGTYPESnumeric_to_int**

Convert a variable of type numeric to int.

```c
int PGTYPESnumeric_to_int(numeric *nv, int *ip);
```

The function converts the numeric value from the variable that nv points to into the integer variable that ip points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable errno will be set to PGTYPES_NUM_OVERFLOW additionally.

**PGTYPESnumeric_to_long**

Convert a variable of type numeric to long.

```c
int PGTYPESnumeric_to_long(numeric *nv, long *lp);
```

The function converts the numeric value from the variable that nv points to into the long integer variable that lp points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable errno will be set to PGTYPES_NUM_OVERFLOW additionally.

**PGTYPESnumeric_to_decimal**

Convert a variable of type numeric to decimal.

```c
int PGTYPESnumeric_to_decimal(numeric *src, decimal *dst);
```

The function converts the numeric value from the variable that src points to into the decimal variable that dst points to. It returns 0 on success and -1 if an error occurs, including overflow. On overflow, the global variable errno will be set to PGTYPES_NUM_OVERFLOW additionally.

**PGTYPESnumeric_from_decimal**

Convert a variable of type decimal to numeric.

```c
int PGTYPESnumeric_from_decimal(decimal *src, numeric *dst);
```

The function converts the decimal value from the variable that src points to into the numeric variable that dst points to. It returns 0 on success and -1 if an error occurs. Since the decimal type is implemented as a limited version of the numeric type, overflow cannot occur with this conversion.

### 34.6.3. The date Type

The date type in C enables your programs to deal with data of the SQL type date. See Section 8.5 for the equivalent type in the PostgreSQL server.

The following functions can be used to work with the date type:

**PGTYPESdate_from_timestamp**

Extract the date part from a timestamp.

```c
date PGTYPESdate_from_timestamp(timestamp dt);
```
The function receives a timestamp as its only argument and returns the extracted date part from this timestamp.

**PGTYPESdate_from_asc**

Parse a date from its textual representation.

```c
date PGTYPESdate_from_asc(char *str, char **endptr);
```

The function receives a C char* string `str` and a pointer to a C char* string `endptr`. At the moment ECPG always parses the complete string and so it currently does not support to store the address of the first invalid character in `*endptr`. You can safely set `endptr` to NULL.

Note that the function always assumes MDY-formatted dates and there is currently no variable to change that within ECPG.

Table 34-2 shows the allowed input formats.

**Table 34-2. Valid Input Formats for PGTYPESdate_from_asc**

<table>
<thead>
<tr>
<th>Input</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>January 8, 1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>1999-01-08</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>1/8/1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>1/18/1999</td>
<td>January 18, 1999</td>
</tr>
<tr>
<td>01/02/03</td>
<td>February 1, 2003</td>
</tr>
<tr>
<td>1999-Jan-08</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>Jan-08-1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>08-Jan-1999</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>99-Jan-08</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>08-Jan-99</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>08-Jan-06</td>
<td>January 8, 2006</td>
</tr>
<tr>
<td>Jan-08-99</td>
<td>January 8, 1999</td>
</tr>
<tr>
<td>19990108</td>
<td>ISO 8601; January 8, 1999</td>
</tr>
<tr>
<td>990108</td>
<td>ISO 8601; January 8, 1999</td>
</tr>
<tr>
<td>1999.008</td>
<td>year and day of year</td>
</tr>
<tr>
<td>J2451187</td>
<td>Julian day</td>
</tr>
<tr>
<td>January 8, 99 BC</td>
<td>year 99 before the Common Era</td>
</tr>
</tbody>
</table>

**PGTYPESdate_to_asc**

Return the textual representation of a date variable.

```c
char *PGTYPESdate_to_asc(date dDate);
```

The function receives the date `dDate` as its only parameter. It will output the date in the form `1999-01-18`, i.e., in the YYYY-MM-DD format. The result must be freed with `PGTYPESchar_free()`.

**PGTYPESdate_julmdy**

Extract the values for the day, the month and the year from a variable of type date.

```c
void PGTYPESdate_julmdy(date d, int *mdy);
```

The function receives the date `d` and a pointer to an array of 3 integer values `mdy`. The variable name indicates the sequential order: `mdy[0]` will be set to contain the number of the month, `mdy[1]` will be set to the value of the day and `mdy[2]` will contain the year.
PGTYPESdate_mdyjul

Create a date value from an array of 3 integers that specify the day, the month and the year of the date.

```c
void PGTYPESdate_mdyjul(int *mdy, date *jdate);
```

The function receives the array of the 3 integers (`mdy`) as its first argument and as its second argument a pointer to a variable of type date that should hold the result of the operation.

PGTYPESdate_dayofweek

Return a number representing the day of the week for a date value.

```c
int PGTYPESdate_dayofweek(date d);
```

The function receives the date variable (`d`) as its only argument and returns an integer that indicates the day of the week for this date.

- 0 - Sunday
- 1 - Monday
- 2 - Tuesday
- 3 - Wednesday
- 4 - Thursday
- 5 - Friday
- 6 - Saturday

PGTYPESdate_today

Get the current date.

```c
void PGTYPESdate_today(date *d);
```

The function receives a pointer to a date variable (`d`) that it sets to the current date.

PGTYPESdate_fmt_asc

Convert a variable of type date to its textual representation using a format mask.

```c
int PGTYPESdate_fmt_asc(date dDate, char *fmtstring, char *outbuf);
```

The function receives the date to convert (`dDate`), the format mask (`fmtstring`) and the string that will hold the textual representation of the date (`outbuf`).

On success, 0 is returned and a negative value if an error occurred.

The following literals are the field specifiers you can use:

- **dd** - The number of the day of the month.
- **mm** - The number of the month of the year.
- **yy** - The number of the year as a two digit number.
- **yyyy** - The number of the year as a four digit number.
- **ddd** - The name of the day (abbreviated).
- **mmm** - The name of the month (abbreviated).

All other characters are copied 1:1 to the output string.

Table 34-3 indicates a few possible formats. This will give you an idea of how to use this function. All output lines are based on the same date: November 23, 1959.
Table 34-3. Valid Input Formats for PGTYPESdate_fmt_asc

<table>
<thead>
<tr>
<th>Format</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>mmddyy</td>
<td>112359</td>
</tr>
<tr>
<td>ddmmyy</td>
<td>231159</td>
</tr>
<tr>
<td>yymmdd</td>
<td>591123</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>59/11/23</td>
</tr>
<tr>
<td>yy mm dd</td>
<td>59 11 23</td>
</tr>
<tr>
<td>yy.mm.dd</td>
<td>59.11.23</td>
</tr>
<tr>
<td>.mm.yyyy.dd.</td>
<td>.11.1959.23.</td>
</tr>
<tr>
<td>mmm. dd, yyyy</td>
<td>Nov. 23, 1959</td>
</tr>
<tr>
<td>mmm dd yyyy</td>
<td>Nov 23 1959</td>
</tr>
<tr>
<td>yyyy dd mm</td>
<td>1959 23 11</td>
</tr>
<tr>
<td>ddd, mmm. dd, yyyy</td>
<td>Mon, Nov. 23, 1959</td>
</tr>
<tr>
<td>(ddd) mmm. dd, yyyy</td>
<td>(Mon) Nov. 23, 1959</td>
</tr>
</tbody>
</table>

PGTYPESdate_defmt_asc

Use a format mask to convert a C char* string to a value of type date.

```
int PGTYPESdate_defmt_asc(date *d, char *fmt, char *str);
```

The function receives a pointer to the date value that should hold the result of the operation (d), the format mask to use for parsing the date (fmt) and the C char* string containing the textual representation of the date (str). The textual representation is expected to match the format mask. However you do not need to have a 1:1 mapping of the string to the format mask. The function only analyzes the sequential order and looks for the literals yy or yyyy that indicate the position of the year, mm to indicate the position of the month and dd to indicate the position of the day.

Table 34-4 indicates a few possible formats. This will give you an idea of how to use this function.

Table 34-4. Valid Input Formats for rdefmtdate

<table>
<thead>
<tr>
<th>Format</th>
<th>String</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>ddmmyy</td>
<td>21-2-54</td>
<td>1954-02-21</td>
</tr>
<tr>
<td>ddmmyy</td>
<td>2-12-54</td>
<td>1954-12-02</td>
</tr>
<tr>
<td>ddmmyy</td>
<td>20111954</td>
<td>1954-11-20</td>
</tr>
<tr>
<td>ddmmyy</td>
<td>2040464</td>
<td>1964-04-13</td>
</tr>
<tr>
<td>mmm.dd.yyyy</td>
<td>MAR-12-1967</td>
<td>1967-03-12</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>1954, February 3rd</td>
<td>1954-02-03</td>
</tr>
<tr>
<td>mmm.dd.yyyy</td>
<td>041269</td>
<td>1969-04-12</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>In the year 2525, in the month of July, mankind will be alive on the 28th day</td>
<td>2525-07-28</td>
</tr>
<tr>
<td>dd-mm-yyyy</td>
<td>I said on the 28th of July in the year 2525</td>
<td>2525-07-28</td>
</tr>
<tr>
<td>mmm.dd.yyyy</td>
<td>9/14/58</td>
<td>1958-09-14</td>
</tr>
<tr>
<td>yy/mm/dd</td>
<td>47/03/29</td>
<td>1947-03-29</td>
</tr>
<tr>
<td>mmm.dd.yyyy</td>
<td>oct 28 1975</td>
<td>1975-10-28</td>
</tr>
</tbody>
</table>
### 34.6.4. The timestamp Type

The timestamp type in C enables your programs to deal with data of the SQL type timestamp. See Section 8.5 for the equivalent type in the PostgreSQL server.

The following functions can be used to work with the timestamp type:

**PGTYPEStimestamp_from_asc**

Parse a timestamp from its textual representation into a timestamp variable.

```c
timestamp PGTYPEStimestamp_from_asc(char *str, char **endptr);
```

The function receives the string to parse (`str`) and a pointer to a C `char*` (`endptr`). At the moment ECPG always parses the complete string and so it currently does not support to store the address of the first invalid character in `*endptr`. You can safely set `endptr` to NULL.

The function returns the parsed timestamp on success. On error, `PGTYPESInvalidTimestamp` is returned and `errno` is set to `PGTYPES_TS_BAD_TIMESTAMP`. See `PGTYPESInvalidTimestamp` for important notes on this value.

In general, the input string can contain any combination of an allowed date specification, a whitespace character and an allowed time specification. Note that time zones are not supported by ECPG. It can parse them but does not apply any calculation as the PostgreSQL server does for example. Timezone specifiers are silently discarded.

Table 34-5 contains a few examples for input strings.

**Table 34-5. Valid Input Formats for PGTYPEStimestamp_from_asc**

<table>
<thead>
<tr>
<th>Input</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>1999-01-08 04:05:06</td>
<td>1999-01-08 04:05:06</td>
</tr>
<tr>
<td>January 8 04:05:06 1999 PST</td>
<td>1999-01-08 04:05:06</td>
</tr>
<tr>
<td>1999-Jan-08 04:05:06.789-8</td>
<td>1999-01-08 04:05:06.789 (time zone specifier ignored)</td>
</tr>
<tr>
<td>J2451187 04:05-08:00</td>
<td>1999-01-08 04:05:00 (time zone specifier ignored)</td>
</tr>
</tbody>
</table>

**PGTYPEStimestamp_to_asc**

Converts a date to a C `char*` string.

```c
char *PGTYPEStimestamp_to_asc(timestamp tstamp);
```

The function receives the timestamp `tstamp` as its only argument and returns an allocated string that contains the textual representation of the timestamp. The result must be freed with `PGTYPESchar_free()`.

**PGTYPEStimestamp_current**

Retrieve the current timestamp.

```c
void PGTYPEStimestamp_current(timestamp *ts);
```
The function retrieves the current timestamp and saves it into the timestamp variable that `ts` points to.

```c
PGTYPEStimestamp_fmt_asc
```

Convert a timestamp variable to a C char* using a format mask.

```c
int PGTYPEStimestamp_fmt_asc(timestamp *ts, char *output, int str_len, char *fmtstr);
```

The function receives a pointer to the timestamp to convert as its first argument (`ts`), a pointer to the output buffer (`output`), the maximal length that has been allocated for the output buffer (`str_len`) and the format mask to use for the conversion (`fmtstr`).

Upon success, the function returns 0 and a negative value if an error occurred.

You can use the following format specifiers for the format mask. The format specifiers are the same ones that are used in the `strftime` function in libc. Any non-format specifier will be copied into the output buffer.

- `%A` - is replaced by national representation of the full weekday name.
- `%a` - is replaced by national representation of the abbreviated weekday name.
- `%B` - is replaced by national representation of the full month name.
- `%b` - is replaced by national representation of the abbreviated month name.
- `%C` - is replaced by `(year / 100)` as decimal number; single digits are preceded by a zero.
- `%c` - is replaced by national representation of time and date.
- `%d` - is equivalent to `%m/%d/%y`.
- `%D` - is equivalent to `%m/%d/%y`.
- `%e` - is replaced by the day of month as a decimal number (1-31); single digits are preceded by a blank.
- `%f` - is equivalent to `%Y-%m-%d`.
- `%g` - is replaced by the same year as in `%G`, but as a decimal number without century (00-99).
- `%H` - is replaced by the hour (24-hour clock) as a decimal number (00-23).
- `%I` - is replaced by the hour (12-hour clock) as a decimal number (01-12).
- `%j` - is replaced by the day of the year as a decimal number (001-366).
- `%k` - is replaced by the hour (24-hour clock) as a decimal number (0-23); single digits are preceded by a blank.
- `%l` - is replaced by the hour (12-hour clock) as a decimal number (1-12); single digits are preceded by a blank.
- `%M` - is replaced by the minute as a decimal number (00-59).
- `%m` - is replaced by the month as a decimal number (01-12).
Chapter 34. ECPG - Embedded SQL in C

- \%n - is replaced by a newline.
- \%O* - the same as \%E*.
- \%p - is replaced by national representation of either “ante meridiem” or “post meridiem” as appropriate.
- \%R - is equivalent to \%H:\%M.
- \%r - is equivalent to \%I:\%M:\%S \%p.
- \%S - is replaced by the second as a decimal number (00-60).
- \%s - is replaced by the number of seconds since the Epoch, UTC.
- \%T - is equivalent to \%H:\%M:\%S
- \%t - is replaced by a tab.
- \%U - is replaced by the week number of the year (Sunday as the first day of the week) as a decimal number (00-53).
- \%u - is replaced by the weekday (Monday as the first day of the week) as a decimal number (1-7).
- \%V - is replaced by the week number of the year (Monday as the first day of the week) as a decimal number (01-53). If the week containing January 1 has four or more days in the new year, then it is week 1; otherwise it is the last week of the previous year, and the next week is week 1.
- \%v - is equivalent to \%e-%b-%Y.
- \%W - is replaced by the week number of the year (Monday as the first day of the week) as a decimal number (00-53).
- \%w - is replaced by the weekday (Sunday as the first day of the week) as a decimal number (0-6).
- \%X - is replaced by national representation of the time.
- \%x - is replaced by national representation of the date.
- \%Y - is replaced by the year with century as a decimal number.
- \%y - is replaced by the year without century as a decimal number (00-99).
- \%Z - is replaced by the time zone name.
- \%z - is replaced by the time zone offset from UTC; a leading plus sign stands for east of UTC, a minus sign for west of UTC, hours and minutes follow with two digits each and no delimiter between them (common form for RFC 822 date headers).
- \%+ - is replaced by national representation of the date and time.
- \%+- - GNU libc extension. Do not do any padding when performing numerical outputs.
- \$_* - GNU libc extension. Explicitly specify space for padding.
- \%0* - GNU libc extension. Explicitly specify zero for padding.
- %% - is replaced by \%.

**PGTYPEStimestamp_sub**

Subtract one timestamp from another one and save the result in a variable of type interval.

```c
int PGTYPEStimestamp_sub(timestamp *ts1, timestamp *ts2, interval *iv);
```
The function will subtract the timestamp variable that \texttt{ts2} points to from the timestamp variable that \texttt{ts1} points to and will store the result in the interval variable that \texttt{iv} points to.

Upon success, the function returns 0 and a negative value if an error occurred.

\begin{verbatim}
PGTYPEStimestamp_defmt_asc
Parse a timestamp value from its textual representation using a formatting mask.

int PGTYPEStimestamp_defmt_asc(char *str, char *fmt, timestamp *d);
\end{verbatim}

The function receives the textual representation of a timestamp in the variable \texttt{str} as well as the formatting mask to use in the variable \texttt{fmt}. The result will be stored in the variable that \texttt{d} points to.

If the formatting mask \texttt{fmt} is NULL, the function will fall back to the default formatting mask which is \texttt{%Y-%m-%d %H:%M:%S}.

This is the reverse function to \texttt{PGTYPEStimestamp_fmt_asc}. See the documentation there in order to find out about the possible formatting mask entries.

\begin{verbatim}
PGTYPEStimestamp_add_interval
Add an interval variable to a timestamp variable.

int PGTYPEStimestamp_add_interval(timestamp *tin, interval *span, timestamp *tout);
\end{verbatim}

The function receives a pointer to a timestamp variable \texttt{tin} and a pointer to an interval variable \texttt{span}. It adds the interval to the timestamp and saves the resulting timestamp in the variable that \texttt{tout} points to.

Upon success, the function returns 0 and a negative value if an error occurred.

\begin{verbatim}
PGTYPEStimestamp_sub_interval
Subtract an interval variable from a timestamp variable.

int PGTYPEStimestamp_sub_interval(timestamp *tin, interval *span, timestamp *tout);
\end{verbatim}

The function subtracts the interval variable that \texttt{span} points to from the timestamp variable that \texttt{tin} points to and saves the result into the variable that \texttt{tout} points to.

Upon success, the function returns 0 and a negative value if an error occurred.

### 34.6.5. The interval Type

The interval type in C enables your programs to deal with data of the SQL type interval. See Section 8.5 for the equivalent type in the PostgreSQL server.

The following functions can be used to work with the interval type:

\begin{verbatim}
PGTYPESinterval_new
Return a pointer to a newly allocated interval variable.

interval *PGTYPESinterval_new(void);
\end{verbatim}

\begin{verbatim}
PGTYPESinterval_free
Release the memory of a previously allocated interval variable.

void PGTYPESinterval_new(interval *intvl);
\end{verbatim}
Chapter 34. ECPG - Embedded SQL in C

34.6.6. The decimal Type

The decimal type is similar to the numeric type. However it is limited to a maximum precision of 30 significant digits. In contrast to the numeric type which can be created on the heap only, the decimal type can be created either on the stack or on the heap (by means of the functions PGTYPESdecimal_new and PGTYPESdecimal_free). There are a lot of other functions that deal with the decimal type in the Informix compatibility mode described in Section 34.15.

The following functions can be used to work with the decimal type and are not only contained in the libcompat library.

PGTYPESdecimal_new

Request a pointer to a newly allocated decimal variable.

decimal *PGTYPESdecimal_new(void);

PGTYPESdecimal_free

Free a decimal type, release all of its memory.

void PGTYPESdecimal_free(decimal *var);

34.6.7. errno Values of pgtypeslib

PGTYPES_NUM_BAD_NUMERIC

An argument should contain a numeric variable (or point to a numeric variable) but in fact its in-memory representation was invalid.
Chapter 34. ECPG - Embedded SQL in C

34.6.8. Special Constants of pgtypeslib

PGTYPESInvalidTimestamp

A value of type timestamp representing an invalid time stamp. This is returned by the function PGTYPEStimestamp_from_asc on parse error. Note that due to the internal representation of the timestamp data type, PGTYPESInvalidTimestamp is also a valid timestamp at the same time. It is set to 1899-12-31 23:59:59. In order to detect errors, make sure that your application does not only test for PGTYPESInvalidTimestamp but also for errno != 0 after each call to PGTYPEStimestamp_from_asc.
34.7. Using Descriptor Areas

An SQL descriptor area is a more sophisticated method for processing the result of a SELECT, FETCH or a DESCRIBE statement. An SQL descriptor area groups the data of one row of data together with metadata items into one data structure. The metadata is particularly useful when executing dynamic SQL statements, where the nature of the result columns might not be known ahead of time. PostgreSQL provides two ways to use Descriptor Areas: the named SQL Descriptor Areas and the C-structure SQLDAs.

34.7.1. Named SQL Descriptor Areas

A named SQL descriptor area consists of a header, which contains information concerning the entire descriptor, and one or more item descriptor areas, which basically each describe one column in the result row.

Before you can use an SQL descriptor area, you need to allocate one:

EXEC SQL ALLOCATE DESCRIPTOR identifier;

The identifier serves as the “variable name” of the descriptor area. When you don’t need the descriptor anymore, you should deallocate it:

EXEC SQL DEALLOCATE DESCRIPTOR identifier;

To use a descriptor area, specify it as the storage target in an INTO clause, instead of listing host variables:

EXEC SQL FETCH NEXT FROM mycursor INTO SQL DESCRIPTOR mydesc;

If the result set is empty, the Descriptor Area will still contain the metadata from the query, i.e. the field names.

For not yet executed prepared queries, the DESCRIBE statement can be used to get the metadata of the result set:

EXEC SQL BEGIN DECLARE SECTION;
char *sql_stmt = "SELECT * FROM table1";
EXEC SQL END DECLARE SECTION;
EXEC SQL PREPARE stmt1 FROM :sql_stmt;
EXEC SQL DESCRIBE stmt1 INTO SQL DESCRIPTOR mydesc;

Before PostgreSQL 9.0, the SQL keyword was optional, so using DESCRIPTOR and SQL DESCRIPTOR produced named SQL Descriptor Areas. Now it is mandatory, omitting the SQL keyword produces SQLDA Descriptor Areas, see Section 34.7.2.

In DESCRIBE and FETCH statements, the INTO and USING keywords can be used to similarly: they produce the result set and the metadata in a Descriptor Area.
Now how do you get the data out of the descriptor area? You can think of the descriptor area as a structure with named fields. To retrieve the value of a field from the header and store it into a host variable, use the following command:

```sql
EXEC SQL GET DESCRIPTOR name :hostvar = field;
```

Currently, there is only one header field defined: `COUNT`, which tells how many item descriptor areas exist (that is, how many columns are contained in the result). The host variable needs to be of an integer type. To get a field from the item descriptor area, use the following command:

```sql
EXEC SQL GET DESCRIPTOR name VALUE num :hostvar = field;
```

`num` can be a literal integer or a host variable containing an integer. Possible fields are:

- **CARDINALITY (integer)**
  - number of rows in the result set

- **DATA**
  - actual data item (therefore, the data type of this field depends on the query)

- **DATETIME_INTERVAL_CODE (integer)**
  - When `TYPE` is 9, `DATETIME_INTERVAL_CODE` will have a value of 1 for `DATE`, 2 for `TIME`, 3 for `TIMESTAMP`, 4 for `TIME WITH TIME ZONE`, or 5 for `TIMESTAMP WITH TIME ZONE`.

- **DATETIME_INTERVAL_PRECISION (integer)**
  - not implemented

- **INDICATOR (integer)**
  - the indicator (indicating a null value or a value truncation)

- **KEY_MEMBER (integer)**
  - not implemented

- **LENGTH (integer)**
  - length of the datum in characters

- **NAME (string)**
  - name of the column

- **NULLABLE (integer)**
  - not implemented

- **OCTET_LENGTH (integer)**
  - length of the character representation of the datum in bytes

- **PRECISION (integer)**
  - precision (for type numeric)

- **RETURNED_LENGTH (integer)**
  - length of the datum in characters

- **RETURNED_OCTET_LENGTH (integer)**
  - length of the character representation of the datum in bytes
Chapter 34. ECPG - Embedded SQL in C

SCALE (integer)
    scale (for type numeric)

TYPE (integer)
    numeric code of the data type of the column

In EXECUTE, DECLARE and OPEN statements, the effect of the INTO and USING keywords are different. A Descriptor Area can also be manually built to provide the input parameters for a query or a cursor and USING SQL DESCRIPTOR name is the way to pass the input parameters into a parameterized query. The statement to build a named SQL Descriptor Area is below:

EXEC SQL SET DESCRIPTOR name VALUE num field = :hostvar;

PostgreSQL supports retrieving more than one record in one FETCH statement and storing the data in host variables in this case assumes that the variable is an array. E.g.:

EXEC SQL BEGIN DECLARE SECTION;
int id[5];
EXEC SQL END DECLARE SECTION;

EXEC SQL FETCH 5 FROM mycursor INTO SQL DESCRIPTOR mydesc;
EXEC SQL GET DESCRIPTOR mydesc VALUE 1 :id = DATA;

34.7.2. SQLDA Descriptor Areas

An SQLDA Descriptor Area is a C language structure which can be also used to get the result set and the metadata of a query. One structure stores one record from the result set.

EXEC SQL include sqlda.h;
sqlda_t *mysqlda;

EXEC SQL FETCH 3 FROM mycursor INTO DESCRIPTOR mysqlda;

Note that the SQL keyword is omitted. The paragraphs about the use cases of the INTO and USING keywords in Section 34.7.1 also apply here with an addition. In a DESCRIBE statement the DESCRIPTOR keyword can be completely omitted if the INTO keyword is used:

EXEC SQL DESCRIBE prepared_statement INTO mysqlda;

The general flow of a program that uses SQLDA is:

1. Prepare a query, and declare a cursor for it.
2. Declare an SQLDA for the result rows.
3. Declare an SQLDA for the input parameters, and initialize them (memory allocation, parameter settings).
4. Open a cursor with the input SQLDA.
5. Fetch rows from the cursor, and store them into an output SQLDA.
6. Read values from the output SQLDA into the host variables (with conversion if necessary).
7. Close the cursor.
8. Free the memory area allocated for the input SQLDA.

34.7.2.1. SQLDA Data Structure

SQLDA uses three data structure types: sqlda_t, sqlvar_t, and struct sqlname.

Tip: PostgreSQL’s SQLDA has a similar data structure to the one in IBM DB2 Universal Database, so some technical information on DB2’s SQLDA could help understanding PostgreSQL’s one better.

34.7.2.1.1. sqlda_t Structure

The structure type sqlda_t is the type of the actual SQLDA. It holds one record. And two or more sqlda_t structures can be connected in a linked list with the pointer in the desc_next field, thus representing an ordered collection of rows. So, when two or more rows are fetched, the application can read them by following the desc_next pointer in each sqlda_t node.

The definition of sqlda_t is:

```c
struct sqlda_struct {
    char    sqldaid[8];
    long    sqldabc;
    short   sqln;
    short   sqld;
    struct sqlda_struct *desc_next;
    struct sqlvar_struct sqlvar[1];
};

typedef struct sqlda_struct sqlda_t;
```

The meaning of the fields is:

sqldaid

It contains the literal string "SQLDA ".

sqldabc

It contains the size of the allocated space in bytes.

sqln

It contains the number of input parameters for a parameterized query in case it’s passed into OPEN, DECLARE or EXECUTE statements using the USING keyword. In case it’s used as output of SELECT, EXECUTE or FETCH statements, its value is the same as sqld statement

sqld

It contains the number of fields in a result set.
If the query returns more than one record, multiple linked SQLDA structures are returned, and `desc_next` holds a pointer to the next entry in the list.

This is the array of the columns in the result set.

### 34.7.2.1.2. sqlvar_t Structure

The structure type `sqlvar_t` holds a column value and metadata such as type and length. The definition of the type is:

```c
struct sqlvar_struct {
    short sqltype;
    short sqllen;
    char  *sqldata;
    short  *sqlind;
    struct sqlname sqlname;
};

typedef struct sqlvar_struct sqlvar_t;
```

The meaning of the fields is:

- **sqltype**
  
  Contains the type identifier of the field. For values, see `enum ECPGttype` in `ecpgtype.h`.

- **sqllen**
  
  Contains the binary length of the field. e.g. 4 bytes for `ECPGt_int`.

- **sqldata**
  
  Points to the data. The format of the data is described in Section 34.4.4.

- **sqlind**
  
  Points to the null indicator. 0 means not null, -1 means null.

- **sqlname**
  
  The name of the field.

### 34.7.2.1.3. struct sqlname Structure

A `struct sqlname` structure holds a column name. It is used as a member of the `sqlvar_t` structure. The definition of the structure is:

```c
#define NAMEDATALEN 64

struct sqlname {
    short length;
};
```
Chapter 34. ECPG - Embedded SQL in C

```c
char data[NAMEDATALEN];
```

The meaning of the fields is:

- **length**
  - Contains the length of the field name.
- **data**
  - Contains the actual field name.

### 34.7.2.2. Retrieving a Result Set Using an SQLDA

The general steps to retrieve a query result set through an SQLDA are:

1. Declare an `sqlda_t` structure to receive the result set.
2. Execute `FETCH/EXECUTE/DESCRIBE` commands to process a query specifying the declared SQLDA.
3. Check the number of records in the result set by looking at `sqln`, a member of the `sqlda_t` structure.
4. Get the values of each column from `sqlvar[0], sqlvar[1], etc., members of the sqlda_t structure`.
5. Go to next row (sqlda_t structure) by following the `desc_next` pointer, a member of the `sqlda_t` structure.
6. Repeat above as you need.

Here is an example retrieving a result set through an SQLDA.

First, declare a `sqlda_t` structure to receive the result set.

```c
sqlda_t *sqlda1;
```

Next, specify the SQLDA in a command. This is a `FETCH` command example.

```sql
EXEC SQL FETCH NEXT FROM cur1 INTO DESCRIPTOR sqlda1;
```

Run a loop following the linked list to retrieve the rows.

```c
sqlda_t *cur_sqlda;
for (cur_sqlda = sqlda1;
    cur_sqlda != NULL;
    cur_sqlda = cur_sqlda->desc_next)
{
    ...
}```
Inside the loop, run another loop to retrieve each column data (sqlvar_t structure) of the row.

```c
for (i = 0; i < cur_sqlda->sqld; i++)
{
    sqlvar_t v = cur_sqlda->sqlvar[i];
    char *sqldata = v.sqldata;
    short sqllen = v.sqllen;
    ... 
}
```

To get a column value, check the sqltype value, a member of the sqlvar_t structure. Then, switch to an appropriate way, depending on the column type, to copy data from the sqlvar field to a host variable.

```c
char var_buf[1024];
switch (v.sqltype)
{
    case ECPGt_char:
        memset(&var_buf, 0, sizeof(var_buf));
        memcpy(&var_buf, sqldata, (sizeof(var_buf) < sqllen ? sizeof(var_buf) - 1 : sqllen));
        break;
    case ECPGt_int: /* integer */
        memcpy(&intval, sqldata, sqllen);
        snprintf(var_buf, sizeof(var_buf), "%d", intval);
        break;
    ...
}
```

### 34.7.2.3. Passing Query Parameters Using an SQLDA

The general steps to use an SQLDA to pass input parameters to a prepared query are:

1. Create a prepared query (prepared statement)
2. Declare a sqlda_t structure as an input SQLDA.
3. Allocate memory area (as sqlda_t structure) for the input SQLDA.
4. Set (copy) input values in the allocated memory.
5. Open a cursor with specifying the input SQLDA.

Here is an example.

First, create a prepared statement.

```sql
EXEC SQL BEGIN DECLARE SECTION;
char query[1024] = "SELECT d.oid, * FROM pg_database d, pg_stat_database s WHERE d.oid = ? AND s.datid = ?";
EXEC SQL END DECLARE SECTION;
```
Chapter 34. ECPG - Embedded SQL in C

EXEC SQL PREPARE stmt1 FROM :query;

Next, allocate memory for an SQLDA, and set the number of input parameters in sqln, a member variable of the sqlda_t structure. When two or more input parameters are required for the prepared query, the application has to allocate additional memory space which is calculated by (nr. of params - 1) * sizeof(sqlvar_t). The example shown here allocates memory space for two input parameters.

sqlda_t *sqlda2;

sqlda2 = (sqlda_t *) malloc(sizeof(sqlda_t) + sizeof(sqlvar_t));
memset(sqlda2, 0, sizeof(sqlda_t) + sizeof(sqlvar_t));

sqlda2->sqln = 2; /* number of input variables */

After memory allocation, store the parameter values into the sqlvar[] array. (This is same array used for retrieving column values when the SQLDA is receiving a result set.) In this example, the input parameters are "postgres", having a string type, and 1, having an integer type.

sqlda2->sqlvar[0].sqltype = ECPGt_char;
sqlda2->sqlvar[0].sqldata = "postgres";
sqlda2->sqlvar[0].sqllen = 8;

int intval = 1;
sqlda2->sqlvar[1].sqltype = ECPGt_int;
sqlda2->sqlvar[1].sqldata = (char *) &intval;
sqlda2->sqlvar[1].sqllen = sizeof(intval);

By opening a cursor and specifying the SQLDA that was set up beforehand, the input parameters are passed to the prepared statement.

EXEC SQL OPEN cur1 USING DESCRIPTOR sqlda2;

Finally, after using input SQLDAs, the allocated memory space must be freed explicitly, unlike SQLDAs used for receiving query results.

free(sqlda2);

34.7.2.4. A Sample Application Using SQLDA

Here is an example program, which describes how to fetch access statistics of the databases, specified by the input parameters, from the system catalogs.

This application joins two system tables, pg_database and pg_stat_database on the database OID, and also fetches and shows the database statistics which are retrieved by two input parameters (a database postgres, and OID 1).

First, declare an SQLDA for input and an SQLDA for output.

EXEC SQL include sqlda.h;
Next, connect to the database, prepare a statement, and declare a cursor for the prepared statement.

```c
int main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    char query[1024] = "SELECT d.oid,* FROM pg_database d, pg_stat_database s WHERE d.oid=s.datid AND ( d.datname=? OR d.oid=? )";
    EXEC SQL END DECLARE SECTION;
    EXEC SQL CONNECT TO testdb AS con1 USER testuser;
    EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;
    EXEC SQL PREPARE stmt1 FROM :query;
    EXEC SQL DECLARE cur1 CURSOR FOR stmt1;
}
```

Next, put some values in the input SQLDA for the input parameters. Allocate memory for the input SQLDA, and set the number of input parameters to sqln. Store type, value, and value length into sqltype, sqldata, and sqllen in the sqlvar structure.

```c
/* Create SQLDA structure for input parameters. */
sqlda2 = (sqlda_t *) malloc(sizeof(sqlda_t) + sizeof(sqlvar_t));
memset(sqlda2, 0, sizeof(sqlda_t) + sizeof(sqlvar_t));
sqlda2->sqln = 2; /* number of input variables */

sqlda2->sqlvar[0].sqltype = ECPGt_char;
sqlda2->sqlvar[0].sqldata = "postgres";
sqlda2->sqlvar[0].sqllen = 8;

intval = 1;
sqlda2->sqlvar[1].sqltype = ECPGt_int;
sqlda2->sqlvar[1].sqldata = (char *)&intval;
sqlda2->sqlvar[1].sqllen = sizeof(intval);
```

After setting up the input SQLDA, open a cursor with the input SQLDA.

```c
/* Open a cursor with input parameters. */
EXEC SQL OPEN cur1 USING DESCRIPTOR sqlda2;
```

Fetch rows into the output SQLDA from the opened cursor. (Generally, you have to call FETCH repeatedly in the loop, to fetch all rows in the result set.)

```c
while (1)
{
    sqlda_t *cur_sqlda;

    /* Assign descriptor to the cursor */
    EXEC SQL FETCH NEXT FROM cur1 INTO DESCRIPTOR sqlda1;
}
Next, retrieve the fetched records from the SQLDA, by following the linked list of the sqlda_t structure.

```c
for (cur_sqlda = sqlda1 ;
     cur_sqlda != NULL ;
     cur_sqlda = cur_sqlda->desc_next)
{
...
```

Read each columns in the first record. The number of columns is stored in sqld, the actual data of the first column is stored in sqlvar[0], both members of the sqlda_t structure.

```c
/* Print every column in a row. */
for (i = 0; i < sqlda1->sqld; i++)
{
    sqlvar_t v = sqlda1->sqlvar[i];
    char *sqldata = v.sqldata;
    short sqllen  = v.sqllen;

    strncpy(name_buf, v.sqlname.data, v.sqlname.length);
    name_buf[v.sqlname.length] = '\0';
```

Now, the column data is stored in the variable v. Copy every datum into host variables, looking at v.sqltype for the type of the column.

```c
switch (v.sqltype) {
    int intval;
    double doubleval;
    unsigned long long int longlongval;

    case ECPGt_char:
        memset(&var_buf, 0, sizeof(var_buf));
        memcpy(&var_buf, sqldata, (sizeof(var_buf) <= sqllen ? sizeof(var_buf)-1 : sqllen));
        break;

    case ECPGt_int: /* integer */
        memcpy(&intval, sqldata, sqllen);
        snprintf(var_buf, sizeof(var_buf), "%d", intval);
        break;

    ...

    default:
        ...
}
```

```c
printf("%s = %s (type: %d)\n", name_buf, var_buf, v.sqltype);
```

Close the cursor after processing all of records, and disconnect from the database.
Chapter 34. ECPG - Embedded SQL in C

EXEC SQL CLOSE cur1;
EXEC SQL COMMIT;

EXEC SQL DISCONNECT ALL;

The whole program is shown in Example 34-1.

Example 34-1. Example SQLDA Program

```c
#include <stdlib.h>
#include <string.h>
#include <stdio.h>
#include <unistd.h>
EXEC SQL include sqlda.h;

sqlda_t *sqlda1; /* descriptor for output */
sqlda_t *sqlda2; /* descriptor for input */
EXEC SQL WHENEVER NOT FOUND DO BREAK;
EXEC SQL WHENEVER SQLERROR STOP;

int main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    char query[1024] = "SELECT d.oid,* FROM pg_database d, pg_stat_database s WHERE d.oid=
    intval;
    unsigned long long int longlongval;
    EXEC SQL END DECLARE SECTION;

    EXEC SQL CONNECT TO uptimedb AS con1 USER uptime;
    EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;

    EXEC SQL PREPARE stmt1 FROM :query;
    EXEC SQL DECLARE cur1 CURSOR FOR stmt1;

    /* Create a SQLDA structure for an input parameter */
    sqlda2 = (sqlda_t *)malloc(sizeof(sqlda_t) + sizeof(sqlvar_t));
    memset(sqlda2, 0, sizeof(sqlda_t) + sizeof(sqlvar_t));
    sqlda2->sqln = 2; /* a number of input variables */

    sqlda2->sqlvar[0].sqltype = ECPGt_char;
    sqlda2->sqlvar[0].sqldata = "postgres";
    sqlda2->sqlvar[0].sqllen = 8;

    intval = 1;
    sqlda2->sqlvar[1].sqltype = ECPGt_int;
    sqlda2->sqlvar[1].sqldata = (char *) &intval;
    sqlda2->sqlvar[1].sqllen = sizeof(intval);

    /* Open a cursor with input parameters. */
    EXEC SQL OPEN cur1 USING DESCRIPTOR sqlda2;
```
while (1) {
    sqlda_t *cur_sqlda;

    /* Assign descriptor to the cursor */
    EXEC SQL FETCH NEXT FROM cur1 INTO DESCRIPTOR sqlda1;

    for (cur_sqlda = sqlda1; 
        cur_sqlda != NULL; 
        cur_sqlda = cur_sqlda->desc_next) {
        int i;
        char name_buf[1024];
        char var_buf[1024];

        /* Print every column in a row. */
        for (i=0; i<cur_sqlda->sqld; i++) {
            sqlvar_t v = cur_sqlda->sqlvar[i];
            char *sqldata = v.sqldata;
            short sqllen = v.sqllen;

            strncpy(name_buf, v.sqlname.data, v.sqlname.length);
            name_buf[v.sqlname.length] = '\0';

            switch (v.sqltype) {
            case ECPGt_char:
                memset(&var_buf, 0, sizeof(var_buf));
                memcpy(&var_buf, sqldata, (sizeof(var_buf) <= sqllen ? sizeof(var_buf) : sqllen));
                break;

            case ECPGt_int: /* integer */
                memcpy(&intval, sqldata, sqllen);
                snprintf(var_buf, sizeof(var_buf), "%d", intval);
                break;

            case ECPGt_long_long: /* bigint */
                memcpy(&longlongval, sqldata, sqllen);
                snprintf(var_buf, sizeof(var_buf), "%lld", longlongval);
                break;

            default:
                {
                    int i;
                    memset(var_buf, 0, sizeof(var_buf));
                    for (i = 0; i < sqllen; i++) {
                        char tmpbuf[16];
                        snprintf(tmpbuf, sizeof(tmpbuf), "%02x", (unsigned char) sqldata[i]);
                        strncat(var_buf, tmpbuf, sizeof(var_buf));
                    }
                }
                break;
            }
    }
}
```c
printf("%s = %s (type: %d)\n", name_buf, var_buf, v.sqltype);
}

printf("\n");
}

EXEC SQL CLOSE cur1;
EXEC SQL COMMIT;
EXEC SQL DISCONNECT ALL;

return 0;
}

The output of this example should look something like the following (some numbers will vary).

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Value</th>
<th>Type</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>datname</td>
<td>template1</td>
<td>1</td>
</tr>
<tr>
<td>datdba</td>
<td>10</td>
<td>1</td>
</tr>
<tr>
<td>encoding</td>
<td>0</td>
<td>5</td>
</tr>
<tr>
<td>datistemplate</td>
<td>t</td>
<td>1</td>
</tr>
<tr>
<td>datallowconn</td>
<td>t</td>
<td>1</td>
</tr>
<tr>
<td>datconnlimit</td>
<td>-1</td>
<td>5</td>
</tr>
<tr>
<td>datlastsysoid</td>
<td>11510</td>
<td>1</td>
</tr>
<tr>
<td>datfrozenxid</td>
<td>379</td>
<td>1</td>
</tr>
<tr>
<td>dattablespace</td>
<td>1663</td>
<td>1</td>
</tr>
<tr>
<td>datconfig</td>
<td></td>
<td>1</td>
</tr>
<tr>
<td>datacl</td>
<td></td>
<td>1</td>
</tr>
<tr>
<td>datid</td>
<td>1</td>
<td>1</td>
</tr>
<tr>
<td>datname</td>
<td>template1</td>
<td>1</td>
</tr>
<tr>
<td>numbackends</td>
<td>0</td>
<td>5</td>
</tr>
<tr>
<td>xact_commit</td>
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<td>9</td>
</tr>
<tr>
<td>xact_rollback</td>
<td>0</td>
<td>9</td>
</tr>
<tr>
<td>blks_read</td>
<td>130</td>
<td>9</td>
</tr>
<tr>
<td>blks_hit</td>
<td>7341714</td>
<td>9</td>
</tr>
<tr>
<td>tup_returned</td>
<td>38262679</td>
<td>9</td>
</tr>
<tr>
<td>tup_fetched</td>
<td>1836281</td>
<td>9</td>
</tr>
<tr>
<td>tup_inserted</td>
<td>0</td>
<td>9</td>
</tr>
<tr>
<td>tup_updated</td>
<td>0</td>
<td>9</td>
</tr>
<tr>
<td>tup_deleted</td>
<td>0</td>
<td>9</td>
</tr>
<tr>
<td>oid</td>
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<td>1</td>
</tr>
<tr>
<td>datname</td>
<td>postgres</td>
<td>1</td>
</tr>
<tr>
<td>datdba</td>
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<td>1</td>
</tr>
<tr>
<td>encoding</td>
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<td>5</td>
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<tr>
<td>datistemplate</td>
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<td>1</td>
</tr>
<tr>
<td>datallowconn</td>
<td>t</td>
<td>1</td>
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<tr>
<td>datconnlimit</td>
<td>-1</td>
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<tr>
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<td>1</td>
</tr>
<tr>
<td>datfrozenxid</td>
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<td>1</td>
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<td>1</td>
</tr>
<tr>
<td>datconfig</td>
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<td>datacl</td>
<td></td>
<td>1</td>
</tr>
<tr>
<td>datid</td>
<td>11511</td>
<td>1</td>
</tr>
<tr>
<td>datname</td>
<td>postgres</td>
<td>1</td>
</tr>
<tr>
<td>numbackends</td>
<td>0</td>
<td>5</td>
</tr>
<tr>
<td>xact_commit</td>
<td>221069</td>
<td>9</td>
</tr>
</tbody>
</table>
```

860
34.8. Error Handling

This section describes how you can handle exceptional conditions and warnings in an embedded SQL program. There are two nonexclusive facilities for this.

- Callbacks can be configured to handle warning and error conditions using the `WHENEVER` command.
- Detailed information about the error or warning can be obtained from the `sqlca` variable.

34.8.1. Setting Callbacks

One simple method to catch errors and warnings is to set a specific action to be executed whenever a particular condition occurs. In general:

```sql
EXEC SQL WHENEVER condition action;
```

`condition` can be one of the following:

- `SQLERROR`
  The specified action is called whenever an error occurs during the execution of an SQL statement.

- `SQLWARNING`
  The specified action is called whenever a warning occurs during the execution of an SQL statement.

- `NOT FOUND`
  The specified action is called whenever an SQL statement retrieves or affects zero rows. (This condition is not an error, but you might be interested in handling it specially.)

`action` can be one of the following:

- `CONTINUE`
  This effectively means that the condition is ignored. This is the default.
GOTO label
GO TO label

Jump to the specified label (using a C goto statement).

SQLPRINT

Print a message to standard error. This is useful for simple programs or during prototyping. The details of the message cannot be configured.

STOP

Call exit(1), which will terminate the program.

DO BREAK

Execute the C statement break. This should only be used in loops or switch statements.

CALL name (args)
DO name (args)

Call the specified C functions with the specified arguments.

The SQL standard only provides for the actions CONTINUE and GOTO (and GO TO).

Here is an example that you might want to use in a simple program. It prints a simple message when a warning occurs and aborts the program when an error happens:

EXEC SQL WHENEVER SQLWARNING SQLPRINT;
EXEC SQL WHENEVER SQLERROR STOP;

The statement EXEC SQL WHENEVER is a directive of the SQL preprocessor, not a C statement. The error or warning actions that it sets apply to all embedded SQL statements that appear below the point where the handler is set, unless a different action was set for the same condition between the first EXEC SQL WHENEVER and the SQL statement causing the condition, regardless of the flow of control in the C program. So neither of the two following C program excerpts will have the desired effect:

/*
 * WRONG
 */
int main(int argc, char *argv[])
{
 ...
 if (verbose) {
   EXEC SQL WHENEVER SQLWARNING SQLPRINT;
 }
 ...
 EXEC SQL SELECT ...;
 ...
}

/*
 * WRONG
 */
int main(int argc, char *argv[])
{
 ...
 set_error_handler();
 ...
 EXEC SQL SELECT ...;
34.8.2. sqlca

For more powerful error handling, the embedded SQL interface provides a global variable with the
name sqlca (SQL communication area) that has the following structure:

```c
struct
{
    char sqlcaid[8];
    long sqlabc;
    long sqlcode;
    struct
    {
        int sqlerrml;
        char sqlerrmc[SQLERRMC_LEN];
    } sqlerrm;
    char sqlerrp[8];
    long sqlerrd[6];
    char sqlwarn[8];
    char sqlstate[5];
} sqlca;
```

(In a multithreaded program, every thread automatically gets its own copy of sqlca. This works
similarly to the handling of the standard C global variable errno.)

sqlca covers both warnings and errors. If multiple warnings or errors occur during the execution of
a statement, then sqlca will only contain information about the last one.

If no error occurred in the last SQL statement, sqlca.sqlcode will be 0 and sqlca.sqlstate
will be "00000". If a warning or error occurred, then sqlca.sqlcode will be negative and
sqlca.sqlstate will be different from "00000". A positive sqlca.sqlcode indicates a
harmless condition, such as that the last query returned zero rows. sqlcode and sqlstate are two
different error code schemes; details appear below.

If the last SQL statement was successful, then sqlca.sqlerrd[1] contains the OID of the processed
row, if applicable, and sqlca.sqlerrd[2] contains the number of processed or returned rows, if
applicable to the command.

In case of an error or warning, sqlca.sqlerrm.sqlerrmc will contain a string that describes the
error. The field sqlca.sqlerrm.sqlerrml contains the length of the error message that is stored
in sqlca.sqlerrm.sqlerrmc (the result of strlen(), not really interesting for a C programmer).
Note that some messages are too long to fit in the fixed-size sqlerrmc array; they will be truncated.

In case of a warning, sqlca.sqlwarn[2] is set to W. (In all other cases, it is set to something different
from W.) If sqlca.sqlwarn[1] is set to W, then a value was truncated when it was stored in a host
variable. sqlca.sqlwarn[0] is set to W if any of the other elements are set to indicate a warning.
The fields sqlcaid, sqlabc, sqlerrp, and the remaining elements of sqlerrd and sqlwarn currently contain no useful information.

The structure sqlca is not defined in the SQL standard, but is implemented in several other SQL database systems. The definitions are similar at the core, but if you want to write portable applications, then you should investigate the different implementations carefully.

Here is one example that combines the use of WHENEVER and sqlca, printing out the contents of sqlca when an error occurs. This is perhaps useful for debugging or prototyping applications, before installing a more “user-friendly” error handler.

EXEC SQL WHENEVER SQLERROR CALL print_sqlca();

```c
void
print_sqlca()
{
    fprintf(stderr, "==== sqlca ====
    fprintf(stderr, "sqlcode: %ld\n", sqlca.sqlcode);
    fprintf(stderr, "sqlerrm.sqlerrml: %d\n", sqlca.sqlerrm.sqlerrml);
    fprintf(stderr, "sqlerrm.sqlerrmc: %s\n", sqlca.sqlerrm.sqlerrmc);
    fprintf(stderr, "sqlerrd: %ld %ld %ld %ld %ld %ld\n", sqlca.sqlerrd[0], sqlca.sqlerrd[1], sqlca.sqlerrd[2], sqlca.sqlerrd[3], sqlca.sqlerrd[4], sqlca.sqlerrd[5]);
    fprintf(stderr, "sqlwarn: %d %d %d %d %d %d %d %d\n", sqlca.sqlwarn[0], sqlca.sqlwarn[1], sqlca.sqlwarn[2], sqlca.sqlwarn[3], sqlca.sqlwarn[4], sqlca.sqlwarn[5], sqlca.sqlwarn[6], sqlca.sqlwarn[7]);
    fprintf(stderr, "sqlstate: %5s\n", sqlca.sqlstate);
    fprintf(stderr, "===============\n    

The result could look as follows (here an error due to a misspelled table name):

==== sqlca ====
sqlcode: -400
sqlerrm.sqlerrml: 49
sqlerrm.sqlerrmc: relation "pg_databasep" does not exist on line 38
sqlerrd: 0 0 0 0 0 0
sqlwarn: 0 0 0 0 0 0 0 0
sqlstate: 42P01
===============
```

### 34.8.3. SQLSTATE VS. SQLCODE

The fields sqlca.sqlstate and sqlca.sqlcode are two different schemes that provide error codes. Both are derived from the SQL standard, but SQLCODE has been marked deprecated in the SQL-92 edition of the standard and has been dropped in later editions. Therefore, new applications are strongly encouraged to use SQLSTATE.

SQLSTATE is a five-character array. The five characters contain digits or upper-case letters that represent codes of various error and warning conditions. SQLSTATE has a hierarchical scheme: the first two characters indicate the general class of the condition, the last three characters indicate a subclass of the general condition. A successful state is indicated by the code 00000. The SQLSTATE codes are for the most part defined in the SQL standard. The PostgreSQL server natively supports SQLSTATE error codes; therefore a high degree of consistency can be achieved by using this error code scheme throughout all applications. For further information see Appendix A.
SQLCODE, the deprecated error code scheme, is a simple integer. A value of 0 indicates success, a positive value indicates success with additional information, a negative value indicates an error. The SQL standard only defines the positive value +100, which indicates that the last command returned or affected zero rows, and no specific negative values. Therefore, this scheme can only achieve poor portability and does not have a hierarchical code assignment. Historically, the embedded SQL processor for PostgreSQL has assigned some specific SQLCODE values for its use, which are listed below with their numeric value and their symbolic name. Remember that these are not portable to other SQL implementations. To simplify the porting of applications to the SQLSTATE scheme, the corresponding SQLSTATE is also listed. There is, however, no one-to-one or one-to-many mapping between the two schemes (indeed it is many-to-many), so you should consult the global SQLSTATE listing in Appendix A in each case.

These are the assigned SQLCODE values:

0 (ECPG_NO_ERROR)
Indicates no error. (SQLSTATE 00000)

100 (ECPG_NOT_FOUND)
This is a harmless condition indicating that the last command retrieved or processed zero rows, or that you are at the end of the cursor. (SQLSTATE 02000)

When processing a cursor in a loop, you could use this code as a way to detect when to abort the loop, like this:

```c
while (1)
{
    EXEC SQL FETCH ... ;
    if (sqlca.sqlcode == ECPG_NOT_FOUND)
        break;
}
```

But WHENEVER NOT FOUND DO BREAK effectively does this internally, so there is usually no advantage in writing this out explicitly.

-12 (ECPG_OUT_OF_MEMORY)
Indicates that your virtual memory is exhausted. The numeric value is defined as -ENOMEM. (SQLSTATE YE001)

-200 (ECPG_UNSUPPORTED)
Indicates the preprocessor has generated something that the library does not know about. Perhaps you are running incompatible versions of the preprocessor and the library. (SQLSTATE YE002)

-201 (ECPG_TOO_MANY_ARGUMENTS)
This means that the command specified more host variables than the command expected. (SQLSTATE 07001 or 07002)

-202 (ECPG_TOO_FEW_ARGUMENTS)
This means that the command specified fewer host variables than the command expected. (SQLSTATE 07001 or 07002)

-203 (ECPG_TOO_MANY_MATCHES)
This means a query has returned multiple rows but the statement was only prepared to store one result row (for example, because the specified variables are not arrays). (SQLSTATE 21000)
Chapter 34. ECPG - Embedded SQL in C

-204 (ECPG_INT_FORMAT)
The host variable is of type \texttt{int} and the datum in the database is of a different type and contains a value that cannot be interpreted as an \texttt{int}. The library uses \texttt{strtol()} for this conversion. (SQLSTATE 42804)

-205 (ECPG_UINT_FORMAT)
The host variable is of type \texttt{unsigned int} and the datum in the database is of a different type and contains a value that cannot be interpreted as an \texttt{unsigned int}. The library uses \texttt{strtoul()} for this conversion. (SQLSTATE 42804)

-206 (ECPG_FLOAT_FORMAT)
The host variable is of type \texttt{float} and the datum in the database is of another type and contains a value that cannot be interpreted as a \texttt{float}. The library uses \texttt{strtof()} for this conversion. (SQLSTATE 42804)

-207 (ECPG_NUMERIC_FORMAT)
The host variable is of type \texttt{numeric} and the datum in the database is of another type and contains a value that cannot be interpreted as a \texttt{numeric} value. (SQLSTATE 42804)

-208 (ECPG_INTERVAL_FORMAT)
The host variable is of type \texttt{interval} and the datum in the database is of another type and contains a value that cannot be interpreted as an \texttt{interval} value. (SQLSTATE 42804)

-209 (ECPG_DATE_FORMAT)
The host variable is of type \texttt{date} and the datum in the database is of another type and contains a value that cannot be interpreted as a \texttt{date} value. (SQLSTATE 42804)

-210 (ECPG_TIMESTAMP_FORMAT)
The host variable is of type \texttt{timestamp} and the datum in the database is of another type and contains a value that cannot be interpreted as a \texttt{timestamp} value. (SQLSTATE 42804)

-211 (ECPG_CONVERT_BOOL)
This means the host variable is of type \texttt{bool} and the datum in the database is neither \texttt{t} nor \texttt{f}. (SQLSTATE 42804)

-212 (ECPG_EMPTY)
The statement sent to the PostgreSQL server was empty. (This cannot normally happen in an embedded SQL program, so it might point to an internal error.) (SQLSTATE YE002)

-213 (ECPG_MISSING_INDICATOR)
A null value was returned and no null indicator variable was supplied. (SQLSTATE 22002)

-214 (ECPG_NO_ARRAY)
An ordinary variable was used in a place that requires an array. (SQLSTATE 42804)

-215 (ECPG_DATA_NOT_ARRAY)
The database returned an ordinary variable in a place that requires array value. (SQLSTATE 42804)

-220 (ECPG_NO_CONN)
The program tried to access a connection that does not exist. (SQLSTATE 08003)
-221 (ECPG_NOT_CONN)
The program tried to access a connection that does exist but is not open. (This is an internal error.) (SQLSTATE YE002)

-230 (ECPG_INVALID_STMT)
The statement you are trying to use has not been prepared. (SQLSTATE 26000)

-239 (ECPG_INFORMIX_DUPLICATE_KEY)
Duplicate key error, violation of unique constraint (Informix compatibility mode). (SQLSTATE 23505)

-240 (ECPG_UNKNOWN_DESCRIPTOR)
The descriptor specified was not found. The statement you are trying to use has not been prepared. (SQLSTATE 33000)

-241 (ECPG_INVALID_DESCRIPTOR_INDEX)
The descriptor index specified was out of range. (SQLSTATE 07009)

-242 (ECPG_UNKNOWN_DESCRIPTOR_ITEM)
An invalid descriptor item was requested. (This is an internal error.) (SQLSTATE YE002)

-243 (ECPG_VAR_NOT_NUMERIC)
During the execution of a dynamic statement, the database returned a numeric value and the host variable was not numeric. (SQLSTATE 07006)

-244 (ECPG_VAR_NOT_CHAR)
During the execution of a dynamic statement, the database returned a non-numeric value and the host variable was numeric. (SQLSTATE 07006)

-284 (ECPG_INFORMIX_SUBSELECT_NOT_ONE)
A result of the subquery is not single row (Informix compatibility mode). (SQLSTATE 21000)

-400 (ECPG_PGSQL)
Some error caused by the PostgreSQL server. The message contains the error message from the PostgreSQL server.

-401 (ECPG_TRANS)
The PostgreSQL server signaled that we cannot start, commit, or rollback the transaction. (SQLSTATE 08007)

-402 (ECPG_CONNECT)
The connection attempt to the database did not succeed. (SQLSTATE 08001)

-403 (ECPG_DUPLICATE_KEY)
Duplicate key error, violation of unique constraint. (SQLSTATE 23505)

-404 (ECPG_SUBSELECT_NOT_ONE)
A result for the subquery is not single row. (SQLSTATE 21000)

-602 (ECPG_WARNING_UNKNOWN_PORTAL)
An invalid cursor name was specified. (SQLSTATE 34000)

-603 (ECPG_WARNING_IN_TRANSACTION)
Transaction is in progress. (SQLSTATE 25001)
34.9. Preprocessor Directives

Several preprocessor directives are available that modify how the ecpg preprocessor parses and processes a file.

34.9.1. Including Files

To include an external file into your embedded SQL program, use:

EXEC SQL INCLUDE filename;
EXEC SQL INCLUDE <filename>;
EXEC SQL INCLUDE "filename";

The embedded SQL preprocessor will look for a file named filename.h, preprocess it, and include it in the resulting C output. Thus, embedded SQL statements in the included file are handled correctly.

The ecpg preprocessor will search a file at several directories in following order:

- current directory
- /usr/local/include
- PostgreSQL include directory, defined at build time (e.g., /usr/local/pgsql/include)
- /usr/include

But when EXEC SQL INCLUDE "filename" is used, only the current directory is searched.

In each directory, the preprocessor will first look for the file name as given, and if not found will append .h to the file name and try again (unless the specified file name already has that suffix).

Note that EXEC SQL INCLUDE is not the same as:

#include <filename.h>

because this file would not be subject to SQL command preprocessing. Naturally, you can continue to use the C #include directive to include other header files.

**Note:** The include file name is case-sensitive, even though the rest of the EXEC SQL INCLUDE command follows the normal SQL case-sensitivity rules.
34.9.2. The define and undef Directives

Similar to the directive \#define that is known from C, embedded SQL has a similar concept:

EXEC SQL DEFINE name;
EXEC SQL DEFINE name value;

So you can define a name:
EXEC SQL DEFINE HAVE_FEATURE;

And you can also define constants:
EXEC SQL DEFINE MYNUMBER 12;
EXEC SQL DEFINE MYSTRING 'abc';

Use \texttt{undef} to remove a previous definition:
EXEC SQL UNDEF MYNUMBER;

Of course you can continue to use the C versions \#define and \#undef in your embedded SQL program. The difference is where your defined values get evaluated. If you use \texttt{EXEC SQL DEFINE} then the \texttt{ecpg} preprocessor evaluates the defines and substitutes the values. For example if you write:

EXEC SQL DEFINE MYNUMBER 12;
... EXEC SQL UPDATE Tbl SET col = MYNUMBER;

then \texttt{ecpg} will already do the substitution and your C compiler will never see any name or identifier \texttt{MYNUMBER}. Note that you cannot use \#define for a constant that you are going to use in an embedded SQL query because in this case the embedded SQL precompiler is not able to see this declaration.

34.9.3. \texttt{ifdef}, \texttt{ifndef}, else, elif, and endif Directives

You can use the following directives to compile code sections conditionally:

EXEC SQL ifdef name;

Checks a \texttt{name} and processes subsequent lines if \texttt{name} has been created with \texttt{EXEC SQL define name}.

EXEC SQL ifndef name;

Checks a \texttt{name} and processes subsequent lines if \texttt{name} has not been created with \texttt{EXEC SQL define name}.

EXEC SQL else;

Starts processing an alternative section to a section introduced by either \texttt{EXEC SQL ifdef name} or \texttt{EXEC SQL ifndef name}.

EXEC SQL elif name;

Checks \texttt{name} and starts an alternative section if \texttt{name} has been created with \texttt{EXEC SQL define name}.
EXEC SQL endif;

Ends an alternative section.

Example:

EXEC SQL ifndef TZVAR;
EXEC SQL SET TIMEZONE TO 'GMT';
EXEC SQL elif TZNAME;
EXEC SQL SET TIMEZONE TO TZNAME;
EXEC SQL else;
EXEC SQL SET TIMEZONE TO TZVAR;
EXEC SQL endif;

34.10. Processing Embedded SQL Programs

Now that you have an idea how to form embedded SQL C programs, you probably want to know how to compile them. Before compiling you run the file through the embedded SQL C preprocessor, which converts the SQL statements you used to special function calls. After compiling, you must link with a special library that contains the needed functions. These functions fetch information from the arguments, perform the SQL command using the libpq interface, and put the result in the arguments specified for output.

The preprocessor program is called `ecpg` and is included in a normal PostgreSQL installation. Embedded SQL programs are typically named with an extension `.pgc` If you have a program file called `progl.pgc`, you can preprocess it by simply calling:

`ecpg progl.pgc`

This will create a file called `progl.c`. If your input files do not follow the suggested naming pattern, you can specify the output file explicitly using the `-o` option.

The preprocessed file can be compiled normally, for example:

`cc -c progl.c`

The generated C source files include header files from the PostgreSQL installation, so if you installed PostgreSQL in a location that is not searched by default, you have to add an option such as `-I/usr/local/pgsql/include` to the compilation command line.

To link an embedded SQL program, you need to include the `libecpg` library, like so:

`cc -o myprog progl.o prog2.o ... -lecpg`

Again, you might have to add an option like `-L/usr/local/pgsql/lib` to that command line.

You can use `pg_config` or `pkg-config` with package name `libecpg` to get the paths for your installation.

If you manage the build process of a larger project using make, it might be convenient to include the following implicit rule to your makefiles:

```makefile
ECPG = ecpg
```
The complete syntax of the ecpg command is detailed in ecpg.

The ecpg library is thread-safe by default. However, you might need to use some threading command-line options to compile your client code.

### 34.11. Library Functions

The libecpg library primarily contains “hidden” functions that are used to implement the functionality expressed by the embedded SQL commands. But there are some functions that can usefully be called directly. Note that this makes your code unportable.

- **ECPGdebug(int on, FILE *stream)** turns on debug logging if called with the first argument non-zero. Debug logging is done on stream. The log contains all SQL statements with all the input variables inserted, and the results from the PostgreSQL server. This can be very useful when searching for errors in your SQL statements.

  **Note:** On Windows, if the ecpg libraries and an application are compiled with different flags, this function call will crash the application because the internal representation of the FILE pointers differ. Specifically, multithreaded/single-threaded, release/debug, and static/dynamic flags should be the same for the library and all applications using that library.

- **ECPGget_PGconn(const char *connection_name)** returns the library database connection handle identified by the given name. If connection_name is set to NULL, the current connection handle is returned. If no connection handle can be identified, the function returns NULL. The returned connection handle can be used to call any other functions from libpq, if necessary.

  **Note:** It is a bad idea to manipulate database connection handles made from ecpg directly with libpq routines.

- **ECPGtransactionStatus(const char *connection_name)** returns the current transaction status of the given connection identified by connection_name. See Section 32.2 and libpq’s PQtransactionStatus() for details about the returned status codes.

- **ECPGstatus(int lineno, const char* connection_name)** returns true if you are connected to a database and false if not. connection_name can be NULL if a single connection is being used.

### 34.12. Large Objects

Large objects are not directly supported by ECPG, but ECPG application can manipulate large objects through the libpq large object functions, obtaining the necessary PGconn object by calling
the ECPGget_PGconn() function. (However, use of the ECPGget_PGconn() function and touching PGconn objects directly should be done very carefully and ideally not mixed with other ECPG database access calls.)

For more details about the ECPGget_PGconn(), see Section 34.11. For information about the large object function interface, see Chapter 33.

Large object functions have to be called in a transaction block, so when autocommit is off, BEGIN commands have to be issued explicitly.

Example 34-2 shows an example program that illustrates how to create, write, and read a large object in an ECPG application.

Example 34-2. ECPG Program Accessing Large Objects

```c
#include <stdio.h>
#include <stdlib.h>
#include <libpq-fe.h>
#include <libpq/libpq-fs.h>

EXEC SQL WHENEVER SQLERROR STOP;

int main(void)
{
    PGconn *conn;
    Oid loid;
    int fd;
    char buf[256];
    int buflen = 256;
    char buf2[256];
    int rc;

    memset(buf, 1, buflen);

    EXEC SQL CONNECT TO testdb AS con1;
    EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;

    conn = ECPGget_PGconn("con1");
    printf("conn = %p\n", conn);

    /* create */
    loid = lo_create(conn, 0);
    if (loid &lt; 0)
        printf("lo_create() failed: %s", PQerrorMessage(conn));

    printf("loid = %d\n", loid);

    /* write test */
    fd = lo_open(conn, loid, INV_READ|INV_WRITE);
    if (fd &lt; 0)
        printf("lo_open() failed: %s", PQerrorMessage(conn));

    printf("fd = %d\n", fd);

    rc = lo_write(conn, fd, buf, buflen);
    if (rc &lt; 0)
        printf("lo_write() failed\n");
}
```
34.13. C++ Applications

ECPG has some limited support for C++ applications. This section describes some caveats.

The `ecpg` preprocessor takes an input file written in C (or something like C) and embedded SQL commands, converts the embedded SQL commands into C language chunks, and finally generates a `.c` file. The header file declarations of the library functions used by the C language chunks that `ecpg` generates are wrapped in `extern "C" { ... }` blocks when used under C++, so they should work seamlessly in C++.

In general, however, the `ecpg` preprocessor only understands C; it does not handle the special syntax and reserved words of the C++ language. So, some embedded SQL code written in C++ application code that uses complicated features specific to C++ might fail to be preprocessed correctly or might not work as expected.

A safe way to use the embedded SQL code in a C++ application is hiding the ECPG calls in a C module, which the C++ application code calls into to access the database, and linking that together with the rest of the C++ code. See Section 34.13.2 about that.
34.13.1. Scope for Host Variables

The `ecpg` preprocessor understands the scope of variables in C. In the C language, this is rather simple because the scopes of variables is based on their code blocks. In C++, however, the class member variables are referenced in a different code block from the declared position, so the `ecpg` preprocessor will not understand the scope of the class member variables.

For example, in the following case, the `ecpg` preprocessor cannot find any declaration for the variable `dbname` in the `test` method, so an error will occur.

class TestCpp
{
    EXEC SQL BEGIN DECLARE SECTION;
    char dbname[1024];
    EXEC SQL END DECLARE SECTION;

    public:
    TestCpp();
    void test();
    ~TestCpp();
};

TestCpp::TestCpp()
{
    EXEC SQL CONNECT TO testdb1;
    EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;
}

void Test::test()
{
    EXEC SQL SELECT current_database() INTO :dbname;
    printf("current_database = %s\n", dbname);
}

TestCpp::~TestCpp()
{
    EXEC SQL DISCONNECT ALL;
}

This code will result in an error like this:

```
ecpg test_cpp.pgc
```

test_cpp.pgc:28: ERROR: variable "dbname" is not declared

To avoid this scope issue, the `test` method could be modified to use a local variable as intermediate storage. But this approach is only a poor workaround, because it uglifies the code and reduces performance.

void TestCpp::test()
{
    EXEC SQL BEGIN DECLARE SECTION;
    char tmp[1024];
    EXEC SQL END DECLARE SECTION;

    EXEC SQL SELECT current_database() INTO :tmp;
    strncpy(dbname, tmp, sizeof(tmp));
Chapter 34. ECPG - Embedded SQL in C

34.13.2. C++ Application Development with External C Module

If you understand these technical limitations of the ecpg preprocessor in C++, you might come to the conclusion that linking C objects and C++ objects at the link stage to enable C++ applications to use ECPG features could be better than writing some embedded SQL commands in C++ code directly. This section describes a way to separate some embedded SQL commands from C++ application code with a simple example. In this example, the application is implemented in C++, while C and ECPG is used to connect to the PostgreSQL server.

Three kinds of files have to be created: a C file (*.pgc), a header file, and a C++ file:

```c
#include "test_mod.h"
#include <stdio.h>

void db_connect()
{
   EXEC SQL CONNECT TO testdb1;
   EXEC SQL SELECT pg_catalog.set_config('search_path', '', false); EXEC SQL COMMIT;
}

void db_test()
{
   EXEC SQL BEGIN DECLARE SECTION;
   char dbname[1024];
   EXEC SQL END DECLARE SECTION;

   EXEC SQL SELECT current_database() INTO :dbname;
   printf("current_database = %s\n", dbname);
}

void db_disconnect()
{
   EXEC SQL DISCONNECT ALL;
}
```

A sub-routine module to execute SQL commands embedded in C. It is going to be converted into test_mod.c by the preprocessor.

```c
A header file with declarations of the functions in the C module (test_mod.pgc). It is included by test_cpp.cpp. This file has to have an extern "C" block around the declarations, because it will be linked from the C++ module.

#define __cplusplus
```
extern "C" {
    #endif

void db_connect();
void db_test();
void db_disconnect();

#ifdef __cplusplus
    }
#endif

test_cpp.cpp

The main code for the application, including the main routine, and in this example a C++ class.
#include "test_mod.h"

class TestCpp
{
    public:
        TestCpp();
        void test();
        ~TestCpp();
    
    TestCpp::TestCpp()
    {
        db_connect();
    }

    void
    TestCpp::test()
    {
        db_test();
    }

    TestCpp::~TestCpp()
    {
        db_disconnect();
    }

    int
    main(void)
    {
        TestCpp *t = new TestCpp();

        t->test();
        return 0;
    }

To build the application, proceed as follows. Convert test_mod.pgc into test_mod.c by running ecpg, and generate test_mod.o by compiling test_mod.c with the C compiler:

ecpg -o test_mod.c test_mod.pgc
cc -c test_mod.c -o test_mod.o
Next, generate `test_cpp.o` by compiling `test_cpp.cpp` with the C++ compiler:

```c
 c++ -c test_cpp.cpp -o test_cpp.o
```

Finally, link these object files, `test_cpp.o` and `test_mod.o`, into one executable, using the C++ compiler driver:

```c
 c++ test_cpp.o test_mod.o -lecpg -o test_cpp
```

### 34.14. Embedded SQL Commands

This section describes all SQL commands that are specific to embedded SQL. Also refer to the SQL commands listed in Reference I, *SQL Commands*, which can also be used in embedded SQL, unless stated otherwise.

**ALLOCATE DESCRIPTOR**

**Name**

`ALLOCATE DESCRIPTOR` — allocate an SQL descriptor area

**Synopsis**

`ALLOCATE DESCRIPTOR name`

**Description**

`ALLOCATE DESCRIPTOR` allocates a new named SQL descriptor area, which can be used to exchange data between the PostgreSQL server and the host program.

Descriptor areas should be freed after use using the `DEALLOCATE DESCRIPTOR` command.

**Parameters**

`name`

A name of SQL descriptor, case sensitive. This can be an SQL identifier or a host variable.

**Examples**

```sql
EXEC SQL ALLOCATE DESCRIPTOR mydesc;
```
Compatibility

ALLOCATE DESCRIPTOR is specified in the SQL standard.

See Also

DEALLOCATE DESCRIPTOR, GET DESCRIPTOR, SET DESCRIPTOR
**CONNECT**

**Name**

CONNECT — establish a database connection

**Synopsis**

CONNECT TO connection_target [ AS connection_name ] [ USER connection_user ]
CONNECT TO DEFAULT
CONNECT connection_user
DATABASE connection_target

**Description**

The CONNECT command establishes a connection between the client and the PostgreSQL server.

**Parameters**

connection_target

connection_target specifies the target server of the connection on one of several forms.

[ database_name@[host][:port] ]

- Connect over TCP/IP
- unix:postgresql://host[:port]/[database_name][?connection_option]
- Connect over Unix-domain sockets
- tcp:postgresql://host[:port]/[database_name][?connection_option]

- Connect over TCP/IP
- SQL string constant
  - containing a value in one of the above forms
- host variable
  - host variable of type char[] or VARCHAR[] containing a value in one of the above forms

connection_name

An optional identifier for the connection, so that it can be referred to in other commands. This can be an SQL identifier or a host variable.

connection_user

The user name for the database connection.

This parameter can also specify user name and password, using one the forms

user_name/password, user_name IDENTIFIED BY password, or user_name USING password.
User name and password can be SQL identifiers, string constants, or host variables.

Use all default connection parameters, as defined by libpq.

Examples

Here a several variants for specifying connection parameters:

EXEC SQL CONNECT TO "connectdb" AS main;
EXEC SQL CONNECT TO "connectdb" AS second;
EXEC SQL CONNECT TO "unix:postgresql://200.46.204.71/connectdb" AS main USER connectuser;
EXEC SQL CONNECT TO "unix:postgresql://localhost/connectdb" AS main USER connectuser;
EXEC SQL CONNECT TO 'connectdb' AS main;
EXEC SQL CONNECT TO 'unix:postgresql://localhost/connectdb' AS main USER :user;
EXEC SQL CONNECT TO :db AS :id;
EXEC SQL CONNECT TO :db USER connectuser USING :pw;
EXEC SQL CONNECT TO @localhost AS main USER connectdb;
EXEC SQL CONNECT TO REGRESSDB1 as main;
EXEC SQL CONNECT TO AS main USER connectdb;
EXEC SQL CONNECT TO connectdb AS :id;
EXEC SQL CONNECT TO connectdb AS main USER connectuser/connectdb;
EXEC SQL CONNECT TO connectdb AS main;
EXEC SQL CONNECT TO connectdb@localhost AS main;
EXEC SQL CONNECT TO tcp:postgresql://localhost:5432/testdb;
EXEC SQL CONNECT TO tcp:postgresql://localhost/connectdb USER connectuser IDENTIFIED BY connectpw;
EXEC SQL CONNECT TO unix:postgresql://localhost/ AS main USER connectdb;
EXEC SQL CONNECT TO unix:postgresql://localhost/connectdb AS main USER connectuser;
EXEC SQL CONNECT TO unix:postgresql://localhost/connectdb USER connectuser IDENTIFIED BY "connectpw";
EXEC SQL CONNECT TO unix:postgresql://localhost/connectdb USER connectuser USING "connectpw";

Here is an example program that illustrates the use of host variables to specify connection parameters:

```c
int
main(void)
{
    EXEC SQL BEGIN DECLARE SECTION;
    char *dbname = "testdb"; /* database name */
    char *user = "testuser"; /* connection user name */
    char *connection = "tcp:postgresql://localhost:5432/testdb"; /* connection string */
    char ver[256]; /* buffer to store the version string */
    EXEC SQL END DECLARE SECTION;

    ECPGdebug(1, stderr);

    EXEC SQL CONNECT TO :dbname USER :user;
    EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;
    EXEC SQL SELECT version() INTO :ver;
    EXEC SQL DISCONNECT;

    printf("version: %s\n", ver);
}
```
EXEC SQL CONNECT TO :connection USER :user;
EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;
EXEC SQL SELECT version() INTO :ver;
EXEC SQL DISCONNECT;

printf("version: %s\n", ver);

return 0;
}

Compatibility

CONNECT is specified in the SQL standard, but the format of the connection parameters is implementation-specific.

See Also

DISCONNECT, SET CONNECTION
DEALLOCATE DESCRIPTOR

Name

DEALLOCATE DESCRIPTOR — deallocate an SQL descriptor area

Synopsis

DEALLOCATE DESCRIPTOR name

Description

DEALLOCATE DESCRIPTOR deallocates a named SQL descriptor area.

Parameters

name

The name of the descriptor which is going to be deallocated. It is case sensitive. This can be an SQL identifier or a host variable.

Examples

EXEC SQL DEALLOCATE DESCRIPTOR mydesc;

Compatibility

DEALLOCATE DESCRIPTOR is specified in the SQL standard.

See Also

ALLOCATE DESCRIPTOR, GET DESCRIPTOR, SET DESCRIPTOR
DECLARE

Name
DEQUEUE — define a cursor

Synopsis
DECLARE cursor_name [ BINARY ] [ INSENSITIVE ] [ [ NO ] SCROLL ] CURSOR [ { WITH | WITHOUT } HOLD ] FOR prepared_name
DECLARE cursor_name [ BINARY ] [ INSENSITIVE ] [ [ NO ] SCROLL ] CURSOR [ { WITH | WITHOUT } HOLD ] FOR query

Description
DEQUEUE declares a cursor for iterating over the result set of a prepared statement. This command has slightly different semantics from the direct SQL command DEQUEUE: Whereas the latter executes a query and prepares the result set for retrieval, this embedded SQL command merely declares a name as a “loop variable” for iterating over the result set of a query; the actual execution happens when the cursor is opened with the OPEN command.

Parameters
cursor_name
A cursor name, case sensitive. This can be an SQL identifier or a host variable.

prepared_name
The name of a prepared query, either as an SQL identifier or a host variable.

query
A SELECT or VALUES command which will provide the rows to be returned by the cursor.

For the meaning of the cursor options, see DEQUEUE.

Examples
Examples declaring a cursor for a query:
EXEC SQL DECLARE C CURSOR FOR SELECT * FROM My_Table;
EXEC SQL DECLARE C CURSOR FOR SELECT Item1 FROM T;
EXEC SQL DECLARE cur1 CURSOR FOR SELECT version();

An example declaring a cursor for a prepared statement:
EXEC SQL PREPARE stmt1 AS SELECT version();
EXEC SQL DECLARE cur1 CURSOR FOR stmt1;
Compatibility

DECLARE is specified in the SQL standard.

See Also

OPEN, CLOSE, DECLARE
**DESCRIBE**

**Name**

DESCRIBE — obtain information about a prepared statement or result set

**Synopsis**

DESCRIBE [ OUTPUT ] prepared_name USING [ SQL ] DESCRIPTOR descriptor_name
DESCRIBE [ OUTPUT ] prepared_name INTO [ SQL ] DESCRIPTOR descriptor_name
DESCRIBE [ OUTPUT ] prepared_name INTO sqlda_name

**Description**

DESCRIBE retrieves metadata information about the result columns contained in a prepared statement, without actually fetching a row.

**Parameters**

prepared_name

The name of a prepared statement. This can be an SQL identifier or a host variable.

descriptor_name

A descriptor name. It is case sensitive. It can be an SQL identifier or a host variable.

sqlda_name

The name of an SQLDA variable.

**Examples**

EXEC SQL ALLOCATE DESCRIPTOR mydesc;
EXEC SQL PREPARE stmt1 FROM :sql_stmt;
EXEC SQL DESCRIBE stmt1 INTO SQL DESCRIPTOR mydesc;
EXEC SQL GET DESCRIPTOR mydesc VALUE 1 :charvar = NAME;
EXEC SQL DEALLOCATE DESCRIPTOR mydesc;

**Compatibility**

DESCRIBE is specified in the SQL standard.

**See Also**

ALLOCATE DESCRIPTOR, GET DESCRIPTOR
**DISCONNECT**

**Name**

DISCONNECT — terminate a database connection

**Synopsis**

DISCONNECT _connection_name_
DISCONNECT [ CURRENT ]
DISCONNECT DEFAULT
DISCONNECT ALL

**Description**

DISCONNECT closes a connection (or all connections) to the database.

**Parameters**

_connection_name_

A database connection name established by the CONNECT command.

CURRENT

Close the “current” connection, which is either the most recently opened connection, or the connection set by the SET CONNECTION command. This is also the default if no argument is given to the DISCONNECT command.

DEFAULT

Close the default connection.

ALL

Close all open connections.

**Examples**

```c
int main(void)
{
  EXEC SQL CONNECT TO testdb AS DEFAULT USER testuser;
  EXEC SQL CONNECT TO testdb AS con1 USER testuser;
  EXEC SQL CONNECT TO testdb AS con2 USER testuser;
  EXEC SQL CONNECT TO testdb AS con3 USER testuser;

  EXEC SQL DISCONNECT CURRENT; /* close con3 */
  EXEC SQL DISCONNECT DEFAULT; /* close DEFAULT */
  EXEC SQL DISCONNECT ALL; /* close con2 and con1 */

  return 0;
}
```
Compatibility

`DISCONNECT` is specified in the SQL standard.

See Also

CONNECT, SET CONNECTION
EXECUTE IMMEDIATE

Name
EXECUTE IMMEDIATE — dynamically prepare and execute a statement

Synopsis
EXECUTE IMMEDIATE string

Description
EXECUTE IMMEDIATE immediately prepares and executes a dynamically specified SQL statement, without retrieving result rows.

Parameters
string
A literal C string or a host variable containing the SQL statement to be executed.

Examples
Here is an example that executes an INSERT statement using EXECUTE IMMEDIATE and a host variable named command:

```c
sprintf(command, "INSERT INTO test (name, amount, letter) VALUES ('db: "ri"', 1, 'f')");
EXEC SQL EXECUTE IMMEDIATE :command;
```

Compatibility
EXECUTE IMMEDIATE is specified in the SQL standard.
GET DESCRIPTOR

Name
GET DESCRIPTOR — get information from an SQL descriptor area

Synopsis
GET DESCRIPTOR descriptor_name :cvariable = descriptor_header_item [, ... ]
GET DESCRIPTOR descriptor_name VALUE column_number :cvariable = descriptor_item [, ... ]

Description
GET DESCRIPTOR retrieves information about a query result set from an SQL descriptor area and stores it into host variables. A descriptor area is typically populated using FETCH or SELECT before using this command to transfer the information into host language variables.

This command has two forms: The first form retrieves descriptor “header” items, which apply to the result set in its entirety. One example is the row count. The second form, which requires the column number as additional parameter, retrieves information about a particular column. Examples are the column name and the actual column value.

Parameters

descriptor_name
A descriptor name.
descriptor_header_item
A token identifying which header information item to retrieve. Only COUNT, to get the number of columns in the result set, is currently supported.
column_number
The number of the column about which information is to be retrieved. The count starts at 1.
descriptor_item
A token identifying which item of information about a column to retrieve. See Section 34.7.1 for a list of supported items.
cvariable
A host variable that will receive the data retrieved from the descriptor area.

Examples
An example to retrieve the number of columns in a result set:

EXEC SQL GET DESCRIPTOR d :d_count = COUNT;
An example to retrieve a data length in the first column:

EXEC SQL GET DESCRIPTOR d VALUE 1 :d_returned_octet_length = RETURNED_OCTET_LENGTH;

An example to retrieve the data body of the second column as a string:

EXEC SQL GET DESCRIPTOR d VALUE 2 :d_data = DATA;

Here is an example for a whole procedure of executing SELECT current_database(); and showing the number of columns, the column data length, and the column data:

```c
int
main(void)
{
  EXEC SQL BEGIN DECLARE SECTION;
  int  d_count;
  char d_data[1024];
  int  d_returned_octet_length;
  EXEC SQL END DECLARE SECTION;

  EXEC SQL CONNECT TO testdb AS con1 USER testuser;
  EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;
  EXEC SQL ALLOCATE DESCRIPTOR d;

  /* Declare, open a cursor, and assign a descriptor to the cursor */
  EXEC SQL DECLARE cur CURSOR FOR SELECT current_database();
  EXEC SQL OPEN cur;
  EXEC SQL FETCH NEXT FROM cur INTO SQL DESCRIPTOR d;

  /* Get a number of total columns */
  EXEC SQL GET DESCRIPTOR d :d_count = COUNT;
  printf("d_count = %d\n", d_count);

  /* Get length of a returned column */
  EXEC SQL GET DESCRIPTOR d VALUE 1 :d_returned_octet_length = RETURNED_OCTET_LENGTH;
  printf("d_returned_octet_length = %d\n", d_returned_octet_length);

  /* Fetch the returned column as a string */
  EXEC SQL GET DESCRIPTOR d VALUE 1 :d_data = DATA;
  printf("d_data = %s\n", d_data);

  /* Closing */
  EXEC SQL CLOSE cur;
  EXEC SQL COMMIT;

  EXEC SQL DEALLOCATE DESCRIPTOR d;
  EXEC SQL DISCONNECT ALL;

  return 0;
}
```

When the example is executed, the result will look like this:

```plaintext
d_count = 1
 d_returned_octet_length = 6
```
d_data = testdb

Compatibility

GET DESCRIPTOR is specified in the SQL standard.

See Also

ALLOCATE DESCRIPTOR, SET DESCRIPTOR
OPEN

Name
OPEN — open a dynamic cursor

Synopsis
OPEN cursor_name
OPEN cursor_name USING value [, ... ]
OPEN cursor_name USING SQL DESCRIPTOR descriptor_name

Description
OPEN opens a cursor and optionally binds actual values to the placeholders in the cursor’s declaration. The cursor must previously have been declared with the DECLARE command. The execution of OPEN causes the query to start executing on the server.

Parameters
cursor_name
The name of the cursor to be opened. This can be an SQL identifier or a host variable.

value
A value to be bound to a placeholder in the cursor. This can be an SQL constant, a host variable, or a host variable with indicator.

descriptor_name
The name of a descriptor containing values to be bound to the placeholders in the cursor. This can be an SQL identifier or a host variable.

Examples
EXEC SQL OPEN a;
EXEC SQL OPEN d USING 1, 'test';
EXEC SQL OPEN c1 USING SQL DESCRIPTOR mydesc;
EXEC SQL OPEN :curname1;

Compatibility
OPEN is specified in the SQL standard.
See Also

DECLARE, CLOSE
**PREPARE**

**Name**

PREPARE — prepare a statement for execution

**Synopsis**

PREPARE name FROM string

**Description**

PREPARE prepares a statement dynamically specified as a string for execution. This is different from the direct SQL statement PREPARE, which can also be used in embedded programs. The EXECUTE command is used to execute either kind of prepared statement.

**Parameters**

*prepared_name*

An identifier for the prepared query.

*string*

A literal C string or a host variable containing a preparable statement, one of the SELECT, INSERT, UPDATE, or DELETE.

**Examples**

char *stmt = "SELECT * FROM test1 WHERE a = ? AND b = ?";

EXEC SQL ALLOCATE DESCRIPTOR outdesc;
EXEC SQL PREPARE foo FROM :stmt;
EXEC SQL EXECUTE foo USING SQL DESCRIPTOR indesc INTO SQL DESCRIPTOR outdesc;

**Compatibility**

PREPARE is specified in the SQL standard.

**See Also**

EXECUTE
SET AUTOCOMMIT

**Name**

SET AUTOCOMMIT — set the autocommit behavior of the current session

**Synopsis**

SET AUTOCOMMIT { = | TO } { ON | OFF }

**Description**

SET AUTOCOMMIT sets the autocommit behavior of the current database session. By default, embedded SQL programs are *not* in autocommit mode, so COMMIT needs to be issued explicitly when desired. This command can change the session to autocommit mode, where each individual statement is committed implicitly.

**Compatibility**

SET AUTOCOMMIT is an extension of PostgreSQL ECPG.
**SET CONNECTION**

**Name**

SET CONNECTION — select a database connection

**Synopsis**

SET CONNECTION [ TO | = ] connection_name

**Description**

SET CONNECTION sets the “current” database connection, which is the one that all commands use unless overridden.

**Parameters**

connection_name

A database connection name established by the CONNECT command.

DEFAULT

Set the connection to the default connection.

**Examples**

EXEC SQL SET CONNECTION TO con2;
EXEC SQL SET CONNECTION = con1;

**Compatibility**

SET CONNECTION is specified in the SQL standard.

**See Also**

CONNECT, DISCONNECT
SET DESCRIPTOR

Name

SET DESCRIPTOR — set information in an SQL descriptor area

Synopsis

SET DESCRIPTOR descriptor_name descriptor_header_item = value [, ... ]
SET DESCRIPTOR descriptor_name VALUE number descriptor_item = value [, ... ]

Description

SET DESCRIPTOR populates an SQL descriptor area with values. The descriptor area is then typically used to bind parameters in a prepared query execution.

This command has two forms: The first form applies to the descriptor “header”, which is independent of a particular datum. The second form assigns values to particular datums, identified by number.

Parameters

descriptor_name

A descriptor name.

descriptor_header_item

A token identifying which header information item to set. Only COUNT, to set the number of descriptor items, is currently supported.

number

The number of the descriptor item to set. The count starts at 1.

descriptor_item

A token identifying which item of information to set in the descriptor. See Section 34.7.1 for a list of supported items.

value

A value to store into the descriptor item. This can be an SQL constant or a host variable.

Examples

EXEC SQL SET DESCRIPTOR indeac COUNT = 1;
EXEC SQL SET DESCRIPTOR indeac VALUE 1 DATA = 2;
EXEC SQL SET DESCRIPTOR indeac VALUE 1 DATA = :val1;
EXEC SQL SET DESCRIPTOR indeac VALUE 2 INDICATOR = :val1, DATA = 'some string';
EXEC SQL SET DESCRIPTOR indeac VALUE 2 INDICATOR = :val2null, DATA = :val2;
Compatibility

`SET DESCRIPTOR` is specified in the SQL standard.

See Also

`ALLOCATE DESCRIPTOR`, `GET DESCRIPTOR`
TYPE

Name
TYPE — define a new data type

Synopsis
TYPE type_name IS ctype

Description
The TYPE command defines a new C type. It is equivalent to putting a typedef into a declare section. This command is only recognized when ecpg is run with the -c option.

Parameters
type_name
The name for the new type. It must be a valid C type name.
ctype
A C type specification.

Examples
EXEC SQL TYPE customer IS
    struct
    {
        varchar name[50];
        int phone;
    };
EXEC SQL TYPE cust_ind IS
    struct ind
    {
        short name_ind;
        short phone_ind;
    };
EXEC SQL TYPE c IS char reference;
EXEC SQL TYPE ind IS union { int integer; short smallint; };
EXEC SQL TYPE intarray IS int[AMOUNT];
EXEC SQL TYPE str IS varchar[BUFFERSIZ];
EXEC SQL TYPE string IS char[11];

Here is an example program that uses EXEC SQL TYPE:

EXEC SQL WHENEVER SQLERROR SQLPRINT;
EXEC SQL TYPE tt IS
  struct
  |
  varchar v[256];
  int i;
};
EXEC SQL TYPE tt_ind IS
  struct ind {
    short v_ind;
    short i_ind;
  };

int main(void)
{
  EXEC SQL BEGIN DECLARE SECTION;
  tt t;
  tt_ind t_ind;
  EXEC SQL END DECLARE SECTION;

  EXEC SQL CONNECT TO testdb AS con1;
  EXEC SQL SELECT pg_catalog.set_config('search_path', ", false); EXEC SQL COMMIT;

  EXEC SQL SELECT current_database(), 256 INTO :t:t_ind LIMIT 1;

  printf("t.v = %s\n", t.v.arr);
  printf("t.i = %d\n", t.i);

  printf("t_ind.v_ind = %d\n", t_ind.v_ind);
  printf("t_ind.i_ind = %d\n", t_ind.i_ind);

  EXEC SQL DISCONNECT con1;

  return 0;
}

The output from this program looks like this:

t.v = testdb
 256
 0

Compatibility
The TYPE command is a PostgreSQL extension.
VAR

Name
VAR — define a variable

Synopsis
VAR varname IS ctype

Description
The VAR command assigns a new C data type to a host variable. The host variable must be previously declared in a declare section.

Parameters

varname
A C variable name.

ctype
A C type specification.

Examples
Exec sql begin declare section;
short a;
exec sql end declare section;
EXEC SQL VAR a IS int;

Compatibility
The VAR command is a PostgreSQL extension.
WHENEVER

Name

WHENEVER — specify the action to be taken when an SQL statement causes a specific class condition to be raised

Synopsis

WHENEVER { NOT FOUND | SQLERROR | SQLWARNING } action

Description

Define a behavior which is called on the special cases (Rows not found, SQL warnings or errors) in the result of SQL execution.

Parameters

See Section 34.8.1 for a description of the parameters.

Examples

EXEC SQL WHENEVER NOT FOUND CONTINUE;
EXEC SQL WHENEVER NOT FOUND DO BREAK;
EXEC SQL WHENEVER SQLWARNING SQLPRINT;
EXEC SQL WHENEVER SQLWARNING DO warn();
EXEC SQL WHENEVER SQLERROR sqlprint;
EXEC SQL WHENEVER SQLERROR CALL print2();
EXEC SQL WHENEVER SQLERROR DO handle_error("select");
EXEC SQL WHENEVER SQLERROR DO sqlnotice(NULL, NONO);
EXEC SQL WHENEVER SQLERROR DO sqlprint();
EXEC SQL WHENEVER SQLERROR GOTO error_label;
EXEC SQL WHENEVER SQLERROR STOP;

A typical application is the use of WHENEVER NOT FOUND BREAK to handle looping through result sets:

int main(void)
{
    EXEC SQL CONNECT TO testdb AS con1;
    EXEC SQL SELECT pg_catalog.set_config(’search_path’, ", false); EXEC SQL COMMIT;
    EXEC SQL ALLOCATE DESCRIPTOR d;
    EXEC SQL DECLARE cur CURSOR FOR SELECT current_database(), ’hoge’, 256;
    EXEC SQL OPEN cur;

    /* when end of result set reached, break out of while loop */
    EXEC SQL WHENEVER NOT FOUND DO BREAK;

    while (1)
WHENEVER

{
    EXEC SQL FETCH NEXT FROM cur INTO SQL DESCRIPTOR d;
    ...
}

EXEC SQL CLOSE cur;
EXEC SQL COMMIT;
EXEC SQL DEALLOCATE DESCRIPTOR d;
EXEC SQL DISCONNECT ALL;

return 0;

}

Compatibility
WHENEVER is specified in the SQL standard, but most of the actions are PostgreSQL extensions.
34.15. Informix Compatibility Mode

epg can be run in a so-called Informix compatibility mode. If this mode is active, it tries to behave as if it were the Informix precompiler for Informix E/SQL. Generally spoken this will allow you to use the dollar sign instead of the EXEC SQL primitive to introduce embedded SQL commands:

```
$int j = 3;
$CONNECT TO :dbname;
$CREATE TABLE test(i INT PRIMARY KEY, j INT);
$INSERT INTO test(i, j) VALUES (?, :j);
$COMMIT;
```

**Note:** There must not be any white space between the $ and a following preprocessor directive, that is, include, define, ifdef, etc. Otherwise, the preprocessor will parse the token as a host variable.

There are two compatibility modes: INFORMIX, INFORMIX_SE

When linking programs that use this compatibility mode, remember to link against libcompat that is shipped with ECPG.

Besides the previously explained syntactic sugar, the Informix compatibility mode ports some functions for input, output and transformation of data as well as embedded SQL statements known from E/SQL to ECPG.

Informix compatibility mode is closely connected to the pgtypeslib library of ECPG. pgtypeslib maps SQL data types to data types within the C host program and most of the additional functions of the Informix compatibility mode allow you to operate on those C host program types. Note however that the extent of the compatibility is limited. It does not try to copy Informix behavior; it allows you to do more or less the same operations and gives you functions that have the same name and the same basic behavior but it is no drop-in replacement if you are using Informix at the moment. Moreover, some of the data types are different. For example, PostgreSQL’s datetime and interval types do not know about ranges like for example YEAR TO MINUTE so you won’t find support in ECPG for that either.

34.15.1. Additional Types

The Informix-special "string" pseudo-type for storing right-trimmed character string data is now supported in Informix-mode without using typedef. In fact, in Informix-mode, ECPG refuses to process source files that contain typedef sometype string;

```
EXEC SQL BEGIN DECLARE SECTION;
string userid; /* this variable will contain trimmed data */
EXEC SQL END DECLARE SECTION;

EXEC SQL FETCH MYCUR INTO :userid;
```
34.15.2. Additional/Missing Embedded SQL Statements

CLOSE DATABASE

This statement closes the current connection. In fact, this is a synonym for ECPG’s DISCONNECT CURRENT:

```
$CLOSE DATABASE; /* close the current connection */
EXEC SQL CLOSE DATABASE;
```

FREE cursor_name

Due to the differences how ECPG works compared to Informix’s ESQL/C (i.e. which steps are purely grammar transformations and which steps rely on the underlying run-time library) there is no FREE cursor_name statement in ECPG. This is because in ECPG, DECLARE CURSOR doesn’t translate to a function call into the run-time library that uses to the cursor name. This means that there’s no run-time bookkeeping of SQL cursors in the ECPG run-time library, only in the PostgreSQL server.

FREE statement_name

FREE statement_name is a synonym for DEALLOCATE PREPARE statement_name.

34.15.3. Informix-compatible SQLDA Descriptor Areas

Informix-compatible mode supports a different structure than the one described in Section 34.7.2. See below:

```
struct sqlvar_compat
{
    short sqltype;
    int sqllen;
    char *sqldata;
    short *sqlind;
    char *sqlname;
    char *sqlformat;
    short sqlitype;
    short sqlilen;
    char *sqlidata;
    int sqlxid;
    char *sqltypename;
    int sqlxid;
    short sqltype;
    short sqlownerlen;
    short sqlsourceid;
    char *sqlownername;
    int sqlsourceid;
    char *sqlilongdata;
    int sqlflags;
    void *sqlreserved;
};

struct sqlda_compat
{
    short sqld;
    struct sqlvar_compat *sqlvar;
};
```
The global properties are:

sqld

The number of fields in the SQLDA descriptor.

sqlvar

Pointer to the per-field properties.

desc_name

Unused, filled with zero-bytes.

desc_occ

Size of the allocated structure.

desc_next

Pointer to the next SQLDA structure if the result set contains more than one record.

reserved

Unused pointer, contains NULL. Kept for Informix-compatibility.

The per-field properties are below, they are stored in the sqlvar array:

sqltype

Type of the field. Constants are in sqltypes.h

sqllen

Length of the field data.

sqldata

Pointer to the field data. The pointer is of char * type, the data pointed by it is in a binary format. Example:

int intval;

switch (sqldata->sqlvar[i].sqltype) {
  case SQLINTEGER:
    intval = *(int *)sqldata->sqlvar[i].sqldata;
    break;
  ...
}


Chapter 34. ECPG - Embedded SQL in C

sqlind

Pointer to the NULL indicator. If returned by DESCRIBE or FETCH then it's always a valid pointer. If used as input for EXECUTE ... USING sqlda; then NULL-pointer value means that the value for this field is non-NULL. Otherwise a valid pointer and sqlitype has to be properly set. Example:

```c
if (*((int2 *)sqldata->sqlvar[i].sqlind != 0)
    printf("value is NULL\n");
```

sqlname

Name of the field. 0-terminated string.

sqlformat

Reserved in Informix, value of PQfformat() for the field.

sqlitype

Type of the NULL indicator data. It's always SQLSMINT when returning data from the server. When the SQLDA is used for a parameterized query, the data is treated according to the set type.

sqlilen

Length of the NULL indicator data.

sqlxid

Extended type of the field, result of PQftype().

sqltypename

sqltypelen

sqlownerlen

sqlsourcetype

sqlownername

sqlsourceid

sqlflags

sqlreserved

Unused.

sqlilongdata

It equals to sqldata if sqlilen is larger than 32kB.

Example:

```c
EXEC SQL INCLUDE sqlda.h;

sqlda_t *sqlda; /* This doesn’t need to be under embedded DECLARE SECTION */

EXEC SQL BEGIN DECLARE SECTION;
char *prep_stmt = "select * from table1";
int i;
EXEC SQL END DECLARE SECTION;

...

EXEC SQL PREPARE mystmt FROM :prep_stmt;
EXEC SQL DESCRIBE mystmt INTO sqlda;
printf("# of fields: %d\n", sqlda->sqlid);
```
for (i = 0; i < sqlda->sqlld; i++)
    printf("field %d: ", sqlda->sqlvar[i]->sqlname);

EXEC SQL DECLARE mycursor CURSOR FOR mystmt;
EXEC SQL OPEN mycursor;
EXEC SQL WHENEVER NOT FOUND GOTO out;

while (1)
{
    EXEC SQL FETCH mycursor USING sqlda;
}

EXEC SQL CLOSE mycursor;

free(sqlda); /* The main structure is all to be free(),
   * sqlda and sqlda->sqlvar is in one allocated area */

For more information, see the sqlda.h header and the
 src/interfaces/ecpg/test/compat_informix/sqlda.pgc regression
test.

34.15.4. Additional Functions

decadd
Add two decimal type values.

int decadd(decimal *arg1, decimal *arg2, decimal *sum);
The function receives a pointer to the first operand of type decimal (arg1), a pointer to the
second operand of type decimal (arg2) and a pointer to a value of type decimal that will contain
the sum (sum). On success, the function returns 0. ECPG_INFORMIX_NUM_OVERFLOW is returned
in case of overflow and ECPG_INFORMIX_NUM_UNDERFLOW in case of underflow. -1 is returned
for other failures and errno is set to the respective errno number of the pgtypeslib.

deccmp
Compare two variables of type decimal.

int deccmp(decimal *arg1, decimal *arg2);
The function receives a pointer to the first decimal value (arg1), a pointer to the second decimal
value (arg2) and returns an integer value that indicates which is the bigger value.
- 1, if the value that arg1 points to is bigger than the value that arg2 points to
- -1, if the value that arg1 points to is smaller than the value that arg2 points to
- 0, if the value that arg1 points to and the value that arg2 points to are equal

deccopy
Copy a decimal value.

void deccopy(decimal *src, decimal *target);
The function receives a pointer to the decimal value that should be copied as the first argument
(src) and a pointer to the target structure of type decimal (target) as the second argument.
Chapter 34. ECPG - Embedded SQL in C

deccvasc

Convert a value from its ASCII representation into a decimal type.

int deccvasc(char *cp, int len, decimal *np);

The function receives a pointer to a string that contains the string representation of the number to be converted (cp) as well as its length len. np is a pointer to the decimal value that saves the result of the operation.

Valid formats are for example: -2.794, +3.44, 592.49E07 or -32.84e-4.

The function returns 0 on success. If overflow or underflow occurred, ECPG_INFORMIX_NUM_OVERFLOW or ECPG_INFORMIX_NUM_UNDERFLOW is returned. If the ASCII representation could not be parsed, ECPG_INFORMIX_BAD_NUMERIC is returned or ECPG_INFORMIX_BAD_EXPONENT if this problem occurred while parsing the exponent.

deccvdbl

Convert a value of type double to a value of type decimal.

int deccvdbl(double dbl, decimal *np);

The function receives the variable of type double that should be converted as its first argument (dbl). As the second argument (np), the function receives a pointer to the decimal variable that should hold the result of the operation.

The function returns 0 on success and a negative value if the conversion failed.

deccvint

Convert a value of type int to a value of type decimal.

int deccvint(int in, decimal *np);

The function receives the variable of type int that should be converted as its first argument (in). As the second argument (np), the function receives a pointer to the decimal variable that should hold the result of the operation.

The function returns 0 on success and a negative value if the conversion failed.

deccvlong

Convert a value of type long to a value of type decimal.

int deccvlong(long lng, decimal *np);

The function receives the variable of type long that should be converted as its first argument (lng). As the second argument (np), the function receives a pointer to the decimal variable that should hold the result of the operation.

The function returns 0 on success and a negative value if the conversion failed.

decdiv

Divide two variables of type decimal.

int decdiv(decimal *n1, decimal *n2, decimal *result);

The function receives pointers to the variables that are the first (n1) and the second (n2) operands and calculates n1/n2. result is a pointer to the variable that should hold the result of the operation.

On success, 0 is returned and a negative value if the division fails. If overflow or underflow occurred, the function returns ECPG_INFORMIX_NUM_OVERFLOW or ECPG_INFORMIX_NUM_UNDERFLOW respectively. If an attempt to divide by zero is observed, the function returns ECPG_INFORMIX_DIVIDE_ZERO.
Chapter 34. ECPG - Embedded SQL in C

decmul

Multiply two decimal values.

```c
int decmul(decimal *n1, decimal *n2, decimal *result);
```

The function receives pointers to the variables that are the first (n1) and the second (n2) operands and calculates \( n1 \times n2 \). result is a pointer to the variable that should hold the result of the operation.

On success, 0 is returned and a negative value if the multiplication fails. If overflow or underflow occurred, the function returns ECPG_INFORMIX_NUM_OVERFLOW or ECPG_INFORMIX_NUM_UNDERFLOW respectively.

decsub

Subtract one decimal value from another.

```c
int decsub(decimal *n1, decimal *n2, decimal *result);
```

The function receives pointers to the variables that are the first (n1) and the second (n2) operands and calculates \( n1 - n2 \). result is a pointer to the variable that should hold the result of the operation.

On success, 0 is returned and a negative value if the subtraction fails. If overflow or underflow occurred, the function returns ECPG_INFORMIX_NUM_OVERFLOW or ECPG_INFORMIX_NUM_UNDERFLOW respectively.

dectoasc

Convert a variable of type decimal to its ASCII representation in a C char* string.

```c
int dectoasc(decimal *np, char *cp, int len, int right)
```

The function receives a pointer to a variable of type decimal (np) that it converts to its textual representation. cp is the buffer that should hold the result of the operation. The parameter right specifies, how many digits right of the decimal point should be included in the output. The result will be rounded to this number of decimal digits. Setting right to -1 indicates that all available decimal digits should be included in the output. If the length of the output buffer, which is indicated by len is not sufficient to hold the textual representation including the trailing zero byte, only a single * character is stored in the result and -1 is returned.

The function returns either -1 if the buffer cp was too small or ECPG_INFORMIX_OUT_OF_MEMORY if memory was exhausted.

dectodbl

Convert a variable of type decimal to a double.

```c
int dectodbl(decimal *np, double *dblp);
```

The function receives a pointer to the decimal value to convert (np) and a pointer to the double variable that should hold the result of the operation (dblp).

On success, 0 is returned and a negative value if the conversion failed.

dectoint

Convert a variable to type decimal to an integer.

```c
int dectoint(decimal *np, int *ip);
```

The function receives a pointer to the decimal value to convert (np) and a pointer to the integer variable that should hold the result of the operation (ip).

On success, 0 is returned and a negative value if the conversion failed. If an overflow occurred, ECPG_INFORMIX_NUM_OVERFLOW is returned.
Chapter 34. ECPG - Embedded SQL in C

Note that the ECPG implementation differs from the Informix implementation. Informix limits an integer to the range from -32767 to 32767, while the limits in the ECPG implementation depend on the architecture (-INT_MAX .. INT_MAX).

\texttt{dectolong}

Convert a variable to type decimal to a long integer.

\begin{verbatim}
int dectolong(decimal *np, long *lngp);
\end{verbatim}

The function receives a pointer to the decimal value to convert (\texttt{np}) and a pointer to the long variable that should hold the result of the operation (\texttt{lngp}).

On success, 0 is returned and a negative value if the conversion failed. If an overflow occurred, \texttt{ECPG\_INFORMIX\_NUM\_OVERFLOW} is returned.

Note that the ECPG implementation differs from the Informix implementation. Informix limits a long integer to the range from -2,147,483,647 to 2,147,483,647, while the limits in the ECPG implementation depend on the architecture (-LONG_MAX .. LONG_MAX).

\texttt{rdatestr}

Converts a date to a C char* string.

\begin{verbatim}
int rdatestr(date d, char *str);
\end{verbatim}

The function receives two arguments, the first one is the date to convert (\texttt{d}) and the second one is a pointer to the target string. The output format is always yyyy-mm-dd, so you need to allocate at least 11 bytes (including the zero-byte terminator) for the string.

The function returns 0 on success and a negative value in case of error.

Note that ECPG’s implementation differs from the Informix implementation. In Informix the format can be influenced by setting environment variables. In ECPG however, you cannot change the output format.

\texttt{rstrdate}

Parse the textual representation of a date.

\begin{verbatim}
int rstrdate(char *str, date *d);
\end{verbatim}

The function receives the textual representation of the date to convert (\texttt{str}) and a pointer to a variable of type date (\texttt{d}). This function does not allow you to specify a format mask. It uses the default format mask of Informix which is mm/dd/yyyy. Internally, this function is implemented by means of \texttt{rdefmtdate}. Therefore, \texttt{rstrdate} is not faster and if you have the choice you should opt for \texttt{rdefmtdate} which allows you to specify the format mask explicitly.

The function returns the same values as \texttt{rdefmtdate}.

\texttt{rtoday}

Get the current date.

\begin{verbatim}
void rtoday(date *d);
\end{verbatim}

The function receives a pointer to a date variable (\texttt{d}) that it sets to the current date.

Internally this function uses the \texttt{PGTYPESdate\_today} function.

\texttt{rjulmdy}

Extract the values for the day, the month and the year from a variable of type date.

\begin{verbatim}
int rjulmdy(date d, short mdy[3]);
\end{verbatim}

The function receives the date \texttt{d} and a pointer to an array of 3 short integer values \texttt{mdy}. The variable name indicates the sequential order: \texttt{mdy[0]} will be set to contain the number of the month, \texttt{mdy[1]} will be set to the value of the day and \texttt{mdy[2]} will contain the year.

The function always returns 0 at the moment.
Internally the function uses the `PGTYPESdate_julmdy` function.

**rdefmtdate**

Use a format mask to convert a character string to a value of type date.

```c
int rdefmtdate(date *d, char *fmt, char *str);
```

The function receives a pointer to the date value that should hold the result of the operation (`d`), the format mask to use for parsing the date (`fmt`) and the C char* string containing the textual representation of the date (`str`). The textual representation is expected to match the format mask. However you do not need to have a 1:1 mapping of the string to the format mask. The function only analyzes the sequential order and looks for the literals `yy` or `yyyy` that indicate the position of the year, `mm` to indicate the position of the month and `dd` to indicate the position of the day.

The function returns the following values:

- 0 - The function terminated successfully.
- `ECPG_INFORMIX_ENOSHORTDATE` - The date does not contain delimiters between day, month and year. In this case the input string must be exactly 6 or 8 bytes long but isn’t.
- `ECPG_INFORMIX_ENOTDMY` - The format string did not correctly indicate the sequential order of year, month and day.
- `ECPG_INFORMIX_BAD_DAY` - The input string does not contain a valid day.
- `ECPG_INFORMIX_BAD_MONTH` - The input string does not contain a valid month.
- `ECPG_INFORMIX_BAD_YEAR` - The input string does not contain a valid year.

Internally this function is implemented to use the `PGTYPESdate_defmt_asc` function. See the reference there for a table of example input.

**rfmtdate**

Convert a variable of type date to its textual representation using a format mask.

```c
int rfmtdate(date d, char *fmt, char *str);
```

The function receives the date to convert (`d`), the format mask (`fmt`) and the string that will hold the textual representation of the date (`str`).

On success, 0 is returned and a negative value if an error occurred.

Internally this function uses the `PGTYPESdate_fmt_asc` function, see the reference there for examples.

**rmdyjul**

Create a date value from an array of 3 short integers that specify the day, the month and the year of the date.

```c
int rmdyjul(short mdy[3], date *d);
```

The function receives the array of the 3 short integers (`mdy`) and a pointer to a variable of type date that should hold the result of the operation.

Currently the function returns always 0.

Internally the function is implemented to use the function `PGTYPESdate_mdyjul`.

**rdayofweek**

Return a number representing the day of the week for a date value.

```c
int rdayofweek(date d);
```
Chapter 34. ECPG - Embedded SQL in C

The function receives the date variable \( d \) as its only argument and returns an integer that indicates the day of the week for this date.

- 0 - Sunday
- 1 - Monday
- 2 - Tuesday
- 3 - Wednesday
- 4 - Thursday
- 5 - Friday
- 6 - Saturday

Internally the function is implemented to use the function \texttt{PGTYPESdate\_dayofweek}.

dtc current

Retrieve the current timestamp.

\begin{verbatim}
void dtcurrent(timestamp *ts);
\end{verbatim}

The function retrieves the current timestamp and saves it into the timestamp variable that \( ts \) points to.

dt cvasc

Parses a timestamp from its textual representation into a timestamp variable.

\begin{verbatim}
int dt cvasc(char *str, timestamp *ts);
\end{verbatim}

The function receives the string to parse (\( str \)) and a pointer to the timestamp variable that should hold the result of the operation (\( ts \)).

The function returns 0 on success and a negative value in case of error.

Internally this function uses the \texttt{PGTYPEStimestamp\_from\_asc} function. See the reference there for a table with example inputs.

dtcvfmtasc

Parses a timestamp from its textual representation using a format mask into a timestamp variable.

\begin{verbatim}
dtcvfmtasc(char *inbuf, char *fmtstr, timestamp *dtvalue)
\end{verbatim}

The function receives the string to parse (\( inbuf \)), the format mask to use (\( fmtstr \)) and a pointer to the timestamp variable that should hold the result of the operation (\( dtvalue \)).

This function is implemented by means of the \texttt{PGTYPEStimestamp\_defmt\_asc} function. See the documentation there for a list of format specifiers that can be used.

The function returns 0 on success and a negative value in case of error.

dt sub

Subtract one timestamp from another and return a variable of type interval.

\begin{verbatim}
int dtsub(timestamp *ts1, timestamp *ts2, interval *iv);
\end{verbatim}

The function will subtract the timestamp variable that \( ts2 \) points to from the timestamp variable that \( ts1 \) points to and will store the result in the interval variable that \( iv \) points to.

Upon success, the function returns 0 and a negative value if an error occurred.

dtt oasc

Convert a timestamp variable to a C char* string.

\begin{verbatim}
int dttoasc(timestamp *ts, char *output);
\end{verbatim}
The function receives a pointer to the timestamp variable to convert (ts) and the string that should hold the result of the operation (output). It converts ts to its textual representation according to the SQL standard, which is YYYY-MM-DD HH:MM:SS.

Upon success, the function returns 0 and a negative value if an error occurred.

dttotfmtasc
Convert a timestamp variable to a C char* using a format mask.

int dttotfmtasc(timestamp *ts, char *output, int str_len, char *fmtstr);
The function receives a pointer to the timestamp to convert as its first argument (ts), a pointer to the output buffer (output), the maximal length that has been allocated for the output buffer (str_len) and the format mask to use for the conversion (fmtstr).

Upon success, the function returns 0 and a negative value if an error occurred.

Internally, this function uses the PGTYPEStimestamp_fmt_asc function. See the reference there for information on what format mask specifiers can be used.

intoasc
Convert an interval variable to a C char* string.

int intoasc(interval *i, char *str);
The function receives a pointer to the interval variable to convert (i) and the string that should hold the result of the operation (str). It converts i to its textual representation according to the SQL standard, which is YYYY-MM-DD HH:MM:SS.

Upon success, the function returns 0 and a negative value if an error occurred.

rfmtlong
Convert a long integer value to its textual representation using a format mask.

int rfmtlong(long lng_val, char *fmt, char *outbuf);
The function receives the long value lng_val, the format mask fmt and a pointer to the output buffer outbuf. It converts the long value according to the format mask to its textual representation.

The format mask can be composed of the following format specifying characters:

• * (asterisk) - if this position would be blank otherwise, fill it with an asterisk.
• & (ampersand) - if this position would be blank otherwise, fill it with a zero.
• # - turn leading zeroes into blanks.
• < - left-justify the number in the string.
• , (comma) - group numbers of four or more digits into groups of three digits separated by a comma.
• . (period) - this character separates the whole-number part of the number from the fractional part.
• - (minus) - the minus sign appears if the number is a negative value.
• + (plus) - the plus sign appears if the number is a positive value.
• ( - this replaces the minus sign in front of the negative number. The minus sign will not appear.
• ) - this character replaces the minus and is printed behind the negative value.
• $ - the currency symbol.
rupshift

Convert a string to upper case.

```c
void rupshift(char *str);
```

The function receives a pointer to the string and transforms every lower case character to upper case.

byleng

Return the number of characters in a string without counting trailing blanks.

```c
int byleng(char *str, int len);
```

The function expects a fixed-length string as its first argument (`str`) and its length as its second argument (`len`). It returns the number of significant characters, that is the length of the string without trailing blanks.

ldchar

Copy a fixed-length string into a null-terminated string.

```c
void ldchar(char *src, int len, char *dest);
```

The function receives the fixed-length string to copy (`src`), its length (`len`) and a pointer to the destination memory (`dest`). Note that you need to reserve at least `len+1` bytes for the string that `dest` points to. The function copies at most `len` bytes to the new location (less if the source string has trailing blanks) and adds the null-terminator.

rgetmsg

```c
int rgetmsg(int msgnum, char *s, int maxsize);
```

This function exists but is not implemented at the moment!

rtypalign

```c
int rtypalign(int offset, int type);
```

This function exists but is not implemented at the moment!

rtypmsize

```c
int rtypmsize(int type, int len);
```

This function exists but is not implemented at the moment!

rtypwidth

```c
int rtypwidth(int sqltype, int sqllen);
```

This function exists but is not implemented at the moment!

rsetnull

Set a variable to NULL.

```c
int rsetnull(int t, char *ptr);
```

The function receives an integer that indicates the type of the variable and a pointer to the variable itself that is cast to a C `char` pointer.

The following types exist:

- **CCHARTYPE** - For a variable of type `char` or `char*`
- **CSHORTTYPE** - For a variable of type `short int`
- **CINTTYPE** - For a variable of type `int`
- **CBOOLTYPE** - For a variable of type `boolean`
- **CFLOATTYPE** - For a variable of type `float`
• CLONGTYPE - For a variable of type long
• CDOUBLETYPE - For a variable of type double
• CDECIMALTYPE - For a variable of type decimal
• CDATETYPE - For a variable of type date
• CDTIMETYPE - For a variable of type timestamp

Here is an example of a call to this function:

```c
char c[] = "abc ";
short s = 17;
int i = -74874;

rsetnull(CCHARTYPE, (char *) c);
rsetnull(CSHORTTYPE, (char *) &s);
rsetnull(CINTTYPE, (char *) &i);
```

risnull
Test if a variable is NULL.

```c
int risnull(int t, char *ptr);
```

The function receives the type of the variable to test (t) as well a pointer to this variable (ptr).
Note that the latter needs to be cast to a char*. See the function rsetnull for a list of possible variable types.

Here is an example of how to use this function:

```c
char c[] = "abc ";
short s = 17;
int i = -74874;

risnull(CCHARTYPE, (char *) c);
risnull(CSHORTTYPE, (char *) &s);
risnull(CINTTYPE, (char *) &i);
```

### 34.15.5. Additional Constants

Note that all constants here describe errors and all of them are defined to represent negative values.

In the descriptions of the different constants you can also find the value that the constants represent in the current implementation. However you should not rely on this number. You can however rely on the fact all of them are defined to represent negative values.

ECPG_INFORMIX_NUM_OVERFLOW

Functions return this value if an overflow occurred in a calculation. Internally it is defined as -1200 (the Informix definition).

ECPG_INFORMIX_NUM_UNDERFLOW

Functions return this value if an underflow occurred in a calculation. Internally it is defined as -1201 (the Informix definition).
Chapter 34. ECPG - Embedded SQL in C

ECPG_INFORMIX_DIVIDE_ZERO
Functions return this value if an attempt to divide by zero is observed. Internally it is defined as -1202 (the Informix definition).

ECPG_INFORMIX_BAD_YEAR
Functions return this value if a bad value for a year was found while parsing a date. Internally it is defined as -1204 (the Informix definition).

ECPG_INFORMIX_BAD_MONTH
Functions return this value if a bad value for a month was found while parsing a date. Internally it is defined as -1205 (the Informix definition).

ECPG_INFORMIX_BAD_DAY
Functions return this value if a bad value for a day was found while parsing a date. Internally it is defined as -1206 (the Informix definition).

ECPG_INFORMIX_ENOSHORTDATE
Functions return this value if a parsing routine needs a short date representation but did not get the date string in the right length. Internally it is defined as -1209 (the Informix definition).

ECPG_INFORMIX_DATE_CONVERT
Functions return this value if an error occurred during date formatting. Internally it is defined as -1210 (the Informix definition).

ECPG_INFORMIX_OUT_OF_MEMORY
Functions return this value if memory was exhausted during their operation. Internally it is defined as -1211 (the Informix definition).

ECPG_INFORMIX_ENOTDMY
Functions return this value if a parsing routine was supposed to get a format mask (like mmddyy) but not all fields were listed correctly. Internally it is defined as -1212 (the Informix definition).

ECPG_INFORMIX_BAD_NUMERIC
Functions return this value either if a parsing routine cannot parse the textual representation for a numeric value because it contains errors or if a routine cannot complete a calculation involving numeric variables because at least one of the numeric variables is invalid. Internally it is defined as -1213 (the Informix definition).

ECPG_INFORMIX_BAD_EXPONENT
Functions return this value if a parsing routine cannot parse an exponent. Internally it is defined as -1216 (the Informix definition).

ECPG_INFORMIX_BAD_DATE
Functions return this value if a parsing routine cannot parse a date. Internally it is defined as -1218 (the Informix definition).

ECPG_INFORMIX_EXTRA_CHARS
Functions return this value if a parsing routine is passed extra characters it cannot parse. Internally it is defined as -1264 (the Informix definition).
34.16. Internals

This section explains how ECPG works internally. This information can occasionally be useful to help users understand how to use ECPG.

The first four lines written by `ecpg` to the output are fixed lines. Two are comments and two are include lines necessary to interface to the library. Then the preprocessor reads through the file and writes output. Normally it just echoes everything to the output.

When it sees an `EXEC SQL` statement, it intervenes and changes it. The command starts with `EXEC SQL` and ends with `;`. Everything in between is treated as an SQL statement and parsed for variable substitution.

Variable substitution occurs when a symbol starts with a colon (`:`). The variable with that name is looked up among the variables that were previously declared within a `EXEC SQL DECLARE` section.

The most important function in the library is `ECPGdo`, which takes care of executing most commands. It takes a variable number of arguments. This can easily add up to 50 or so arguments, and we hope this will not be a problem on any platform.

The arguments are:

A line number

This is the line number of the original line; used in error messages only.

A string

This is the SQL command that is to be issued. It is modified by the input variables, i.e., the variables that where not known at compile time but are to be entered in the command. Where the variables should go the string contains `?`.

Input variables

Every input variable causes ten arguments to be created. (See below.)

`ECPGt_EOIT`

An `enum` telling that there are no more input variables.

Output variables

Every output variable causes ten arguments to be created. (See below.) These variables are filled by the function.

`ECPGt_EORT`

An `enum` telling that there are no more variables.

For every variable that is part of the SQL command, the function gets ten arguments:

1. The type as a special symbol.
2. A pointer to the value or a pointer to the pointer.
3. The size of the variable if it is a `char` or `varchar`.
4. The number of elements in the array (for array fetches).
5. The offset to the next element in the array (for array fetches).
6. The type of the indicator variable as a special symbol.
7. A pointer to the indicator variable.
8. 0
9. The number of elements in the indicator array (for array fetches).
10. The offset to the next element in the indicator array (for array fetches).

Note that not all SQL commands are treated in this way. For instance, an open cursor statement like:

```c
EXEC SQL OPEN cursor;
```

is not copied to the output. Instead, the cursor’s DECLARE command is used at the position of the OPEN command because it indeed opens the cursor.

Here is a complete example describing the output of the preprocessor of a file `foo.pgc` (details might change with each particular version of the preprocessor):

```c
EXEC SQL BEGIN DECLARE SECTION;
int index;
int result;
EXEC SQL END DECLARE SECTION;
...
EXEC SQL SELECT res INTO :result FROM mytable WHERE index = :index;
```

is translated into:

```c
/* Processed by ecpg (2.6.0) */
/* These two include files are added by the preprocessor */
#include <ecpgtype.h>
#include <ecpglib.h>

/* exec sql begin declare section */

#line 1 "foo.pgc"
int index;
int result;
/* exec sql end declare section */
...
ECPGdo(__LINE__, NULL, "SELECT res FROM mytable WHERE index = ?",
      ECPGT_int,&(index),1L,1L,sizeof(int),
      ECPGT_NO_INDICATOR, NULL , 0L, 0L, 0L, ECPGT_EOIT,
      ECPGT_int,&(result),1L,1L,sizeof(int),
      ECPGT_NO_INDICATOR, NULL , 0L, 0L, 0L, ECPGT_EORT);
#line 147 "foo.pgc"
```

(The indentation here is added for readability and not something the preprocessor does.)
Chapter 35. The Information Schema

The information schema consists of a set of views that contain information about the objects defined in the current database. The information schema is defined in the SQL standard and can therefore be expected to be portable and remain stable — unlike the system catalogs, which are specific to PostgreSQL and are modeled after implementation concerns. The information schema views do not, however, contain information about PostgreSQL-specific features; to inquire about those you need to query the system catalogs or other PostgreSQL-specific views.

**Note:** When querying the database for constraint information, it is possible for a standard-compliant query that expects to return one row to return several. This is because the SQL standard requires constraint names to be unique within a schema, but PostgreSQL does not enforce this restriction. PostgreSQL automatically-generated constraint names avoid duplicates in the same schema, but users can specify such duplicate names.

This problem can appear when querying information schema views such as `check_constraint_routine_usage`, `check_constraints`, `domain_constraints`, and `referential_constraints`. Some other views have similar issues but contain the table name to help distinguish duplicate rows, e.g., `constraint_column_usage`, `constraint_table_usage`, `table_constraints`.

### 35.1. The Schema

The information schema itself is a schema named `information_schema`. This schema automatically exists in all databases. The owner of this schema is the initial database user in the cluster, and that user naturally has all the privileges on this schema, including the ability to drop it (but the space savings achieved by that are minuscule).

By default, the information schema is not in the schema search path, so you need to access all objects in it through qualified names. Since the names of some of the objects in the information schema are generic names that might occur in user applications, you should be careful if you want to put the information schema in the path.

### 35.2. Data Types

The columns of the information schema views use special data types that are defined in the information schema. These are defined as simple domains over ordinary built-in types. You should not use these types for work outside the information schema, but your applications must be prepared for them if they select from the information schema.

These types are:

- **cardinal_number**
  
  A nonnegative integer.

- **character_data**

  A character string (without specific maximum length).
sql_identifier

A character string. This type is used for SQL identifiers, the type character_data is used for any other kind of text data.

time_stamp

A domain over the type timestamp with time zone

yes_or_no

A character string domain that contains either YES or NO. This is used to represent Boolean (true/false) data in the information schema. (The information schema was invented before the type boolean was added to the SQL standard, so this convention is necessary to keep the information schema backward compatible.)

Every column in the information schema has one of these five types.

### 35.3. information_schema_catalog_name

information_schema_catalog_name is a table that always contains one row and one column containing the name of the current database (current catalog, in SQL terminology).

Table 35-1. information_schema_catalog_name Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>catalog_name</td>
<td>sql_identifier</td>
<td>Name of the database that contains this information schema</td>
</tr>
</tbody>
</table>

### 35.4. administrable_role_authorizations

The view administrable_role_authorizations identifies all roles that the current user has the admin option for.

Table 35-2. administrable_role_authorizations Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role to which this role membership was granted (can be the current user, or a different role in case of nested role memberships)</td>
</tr>
<tr>
<td>role_name</td>
<td>sql_identifier</td>
<td>Name of a role</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>Always YES</td>
</tr>
</tbody>
</table>

### 35.5. applicable_roles

The view applicable_roles identifies all roles whose privileges the current user can use. This
means there is some chain of role grants from the current user to the role in question. The current user itself is also an applicable role. The set of applicable roles is generally used for permission checking.

### Table 35-3. applicable_roles Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role to which this role membership was granted (can be the current user, or a different role in case of nested role memberships)</td>
</tr>
<tr>
<td>role_name</td>
<td>sql_identifier</td>
<td>Name of a role</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the grantee has the admin option on the role, NO if not</td>
</tr>
</tbody>
</table>

#### 35.6. attributes

The view attributes contains information about the attributes of composite data types defined in the database. (Note that the view does not give information about table columns, which are sometimes called attributes in PostgreSQL contexts.) Only those attributes are shown that the current user has access to (by way of being the owner of or having some privilege on the type).

### Table 35-4. attributes Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the data type (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the data type</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the data type</td>
</tr>
<tr>
<td>attribute_name</td>
<td>sql_identifier</td>
<td>Name of the attribute</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the attribute within the data type (count starts at 1)</td>
</tr>
<tr>
<td>attribute_default</td>
<td>character_data</td>
<td>Default expression of the attribute</td>
</tr>
<tr>
<td>is_nullable</td>
<td>yes_or_no</td>
<td>YES if the attribute is possibly nullable, NO if it is known not nullable.</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>----------------------</td>
<td>-------------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the attribute, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in attribute_udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character or bit string type, the declared maximum length; null for all other data types or if no maximum length was declared.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character type, the maximum possible length in octets (bytes) of a datum; null for all other data types. The maximum octet length depends on the declared character maximum length (see above) and the server encoding.</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the attribute (always the current database), null if default or the data type of the attribute is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the attribute, null if default or the data type of the attribute is not collatable</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the attribute, null if default or the data type of the attribute is not collatable</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------</td>
<td>-----------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies a numeric type, this column contains the (declared or implicit) precision of the type for this attribute. The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column <code>numeric_precision_radix</code>. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies a numeric type, this column indicates in which base the values in the columns <code>numeric_precision</code> and <code>numeric_scale</code> are expressed. The value is either 2 or 10. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies an exact numeric type, this column contains the (declared or implicit) scale of the type for this attribute. The scale indicates the number of significant digits to the right of the decimal point. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column <code>numeric_precision_radix</code>. For all other data types, this column is null.</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies a date, time, timestamp, or interval type, this column contains the (declared or implicit) fractional seconds precision of the type for this attribute, that is, the number of decimal digits maintained following the decimal point in the seconds value. For all other data types, this column is null.</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>If data_type identifies an interval type, this column contains the specification which fields the intervals include for this attribute, e.g., YEAR TO MONTH, DAY TO SECOND, etc. If no field restrictions were specified (that is, the interval accepts all fields), and for all other data types, this field is null.</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL (see datetime_precision for the fractional seconds precision of interval type attributes)</td>
</tr>
<tr>
<td>attribute_udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the attribute data type is defined in (always the current database)</td>
</tr>
<tr>
<td>attribute_udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the attribute data type is defined in</td>
</tr>
<tr>
<td>attribute_udt_name</td>
<td>sql_identifier</td>
<td>Name of the attribute data type</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the column, unique among the data type descriptors pertaining to the table. This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
<tr>
<td>is_derived_reference_attribute</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>

See also under Section 35.16, a similarly structured view, for further information on some of the
35.7. character_sets

The view `character_sets` identifies the character sets available in the current database. Since PostgreSQL does not support multiple character sets within one database, this view only shows one, which is the database encoding.

Take note of how the following terms are used in the SQL standard:

character repertoire

An abstract collection of characters, for example `UNICODE`, `UCS`, or `LATIN1`. Not exposed as an SQL object, but visible in this view.

character encoding form

An encoding of some character repertoire. Most older character repertoires only use one encoding form, and so there are no separate names for them (e.g., `LATIN1` is an encoding form applicable to the `LATIN1` repertoire). But for example Unicode has the encoding forms `UTF8`, `UTF16`, etc. (not all supported by PostgreSQL). Encoding forms are not exposed as an SQL object, but are visible in this view.

character set

A named SQL object that identifies a character repertoire, a character encoding, and a default collation. A predefined character set would typically have the same name as an encoding form, but users could define other names. For example, the character set `UTF8` would typically identify the character repertoire `UCS`, encoding form `UTF8`, and some default collation.

You can think of an “encoding” in PostgreSQL either as a character set or a character encoding form. They will have the same name, and there can only be one in one database.

Table 35-5. character_sets Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null.</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null.</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Name of the character set, currently implemented as showing the name of the database encoding</td>
</tr>
<tr>
<td>character_repertoire</td>
<td>sql_identifier</td>
<td>Character repertoire, showing <code>UCS</code> if the encoding is <code>UTF8</code>, else just the encoding name</td>
</tr>
<tr>
<td>form_of_use</td>
<td>sql_identifier</td>
<td>Character encoding form, same as the database encoding</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>default_collate_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the default collation (always the current database, if any collation is identified)</td>
</tr>
<tr>
<td>default_collate_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the default collation</td>
</tr>
<tr>
<td>default_collate_name</td>
<td>sql_identifier</td>
<td>Name of the default collation. The default collation is identified as the collation that matches the COLLATE and CTYPE settings of the current database. If there is no such collation, then this column and the associated schema and catalog columns are null.</td>
</tr>
</tbody>
</table>

#### 35.8. check_constraint_routine_usage

The view `check_constraint_routine_usage` identifies routines (functions and procedures) that are used by a check constraint. Only those routines are shown that are owned by a currently enabled role.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 35.40 for more information.</td>
</tr>
</tbody>
</table>

#### 35.9. check_constraints

The view `check_constraints` contains all check constraints, either defined on a table or on a domain, that are owned by a currently enabled role. (The owner of the table or domain is the owner of the constraint.)
Table 35-7. check_constraints Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>check_clause</td>
<td>character_data</td>
<td>The check expression of the check constraint</td>
</tr>
</tbody>
</table>

35.10. collations

The view collations contains the collations available in the current database.

Table 35-8. collations Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation (always the current database)</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the default collation</td>
</tr>
<tr>
<td>pad_attribute</td>
<td>character_data</td>
<td>Always NO PAD (The alternative PAD SPACE is not supported by PostgreSQL)</td>
</tr>
</tbody>
</table>

35.11. collation_character_set_applicability

The view collation_character_set_applicability identifies which character set the available collations are applicable to. In PostgreSQL, there is only one character set per database (see explanation in Section 35.7), so this view does not provide much useful information.

Table 35-9. collation_character_set_applicability Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation (always the current database)</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the default collation</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Character sets are currently not implemented as schema objects, so this column is null</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Name of the character set</td>
</tr>
</tbody>
</table>

#### 35.12. column_domain_usage

The view `column_domain_usage` identifies all columns (of a table or a view) that make use of some domain defined in the current database and owned by a currently enabled role.

Table 35-10. column_domain_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the domain</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
</tbody>
</table>

#### 35.13. column_options

The view `column_options` contains all the options defined for foreign table columns in the current database. Only those foreign table columns are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-11. column_options Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the foreign table</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the foreign table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

#### 35.14. column_privileges

The view `column_privileges` identifies all privileges granted on columns to a currently enabled role or by a currently enabled role. There is one row for each combination of column, grantor, and grantee.

If a privilege has been granted on an entire table, it will show up in this view as a grant for each column, but only for the privilege types where column granularity is possible: `SELECT`, `INSERT`, `UPDATE`, `REFERENCES`.

**Table 35-12. column_privileges Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Type of the privilege: <code>SELECT</code>, <code>INSERT</code>, <code>UPDATE</code>, or <code>REFERENCES</code></td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

#### 35.15. column_udt_usage

The view `column_udt_usage` identifies all columns that use data types owned by a currently enabled role. Note that in PostgreSQL, built-in data types behave like user-defined types, so they are included here as well. See also Section 35.16 for details.
Table 35-13. column_udt_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the column data type (the underlying type of the domain, if applicable) is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the column data type (the underlying type of the domain, if applicable) is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the column data type (the underlying type of the domain, if applicable)</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
</tbody>
</table>

35.16. columns

The view columns contains information about all table columns (or view columns) in the database. System columns (oid, etc.) are not included. Only those columns are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-14. columns Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the column within the table (count starts at 1)</td>
</tr>
<tr>
<td>column_default</td>
<td>character_data</td>
<td>Default expression of the column</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>-----------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>is_nullable</td>
<td>yes_or_no</td>
<td>YES if the column is possibly nullable, NO if it is known not nullable. A not-null constraint is one way a column can be known not nullable, but there can be others.</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the column, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in udt_name and associated columns). If the column is based on a domain, this column refers to the type underlying the domain (and the domain is identified in domain_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character or bit string type, the declared maximum length; null for all other data types or if no maximum length was declared.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>If data_type identifies a character type, the maximum possible length in octets (bytes) of a datum; null for all other data types. The maximum octet length depends on the declared character maximum length (see above) and the server encoding.</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>If data_type identifies a numeric type, this column contains the (declared or implicit) precision of the type for this column. The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies a numeric type, this column indicates in which base the values in the columns <code>numeric_precision</code> and <code>numeric_scale</code> are expressed. The value is either 2 or 10. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies an exact numeric type, this column contains the (declared or implicit) scale of the type for this column. The scale indicates the number of significant digits to the right of the decimal point. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column <code>numeric_precision_radix</code>. For all other data types, this column is null.</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>If <code>data_type</code> identifies a date, time, timestamp, or interval type, this column contains the (declared or implicit) fractional seconds precision of the type for this column, that is, the number of decimal digits maintained following the decimal point in the seconds value. For all other data types, this column is null.</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>If <code>data_type</code> identifies an interval type, this column contains the specification which fields the intervals include for this column, e.g., YEAR TO MONTH, DAY TO SECOND, etc. If no field restrictions were specified (that is, the interval accepts all fields), and for all other data types, this field is null.</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------------</td>
<td>---------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL (see datetime_precision for the fractional seconds precision of interval type columns)</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the column (always the current database), null if default or the data type of the column is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the column, null if default or the data type of the column is not collatable</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the column, null if default or the data type of the column is not collatable</td>
</tr>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>If the column has a domain type, the name of the database that the domain is defined in (always the current database), else null.</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>If the column has a domain type, the name of the schema that the domain is defined in, else null.</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>If the column has a domain type, the name of the domain, else null.</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the column data type (the underlying type of the domain, if applicable) is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the column data type (the underlying type of the domain, if applicable) is defined in</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------</td>
<td>-------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the column data type (the underlying type of the domain, if applicable)</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the column, unique among the data type descriptors pertaining to the table. This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
<tr>
<td>is_self_referencing</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_identity</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>identity_generation</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>identity_start</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>identity_increment</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>identity_maximum</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>identity_minimum</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>identity_cycle</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>is_generated</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>generation_expression</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>
Since data types can be defined in a variety of ways in SQL, and PostgreSQL contains additional ways to define data types, their representation in the information schema can be somewhat difficult. The column data_type is supposed to identify the underlying built-in type of the column. In PostgreSQL, this means that the type is defined in the system catalog schema pg_catalog. This column might be useful if the application can handle the well-known built-in types specially (for example, format the numeric types differently or use the data in the precision columns). The columns udt_name, udt_schema, and udt_catalog always identify the underlying data type of the column, even if the column is based on a domain. (Since PostgreSQL treats built-in types like user-defined types, built-in types appear here as well. This is an extension of the SQL standard.) These columns should be used if an application wants to process data differently according to the type, because in that case it wouldn’t matter if the column is really based on a domain. If the column is based on a domain, the identity of the domain is stored in the columns domain_name, domain_schema, and domain_catalog. If you want to pair up columns with their associated data types and treat domains as separate types, you could write coalesce(domain_name, udt_name), etc.

### 35.17. `constraint_column_usage`

The view `constraint_column_usage` identifies all columns in the current database that are used by some constraint. Only those columns are shown that are contained in a table owned by a currently enabled role. For a check constraint, this view identifies the columns that are used in the check expression. For a foreign key constraint, this view identifies the columns that the foreign key references. For a unique or primary key constraint, this view identifies the constrained columns.

#### Table 35-15. `constraint_column_usage` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column that is used by some constraint (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column that is used by some constraint</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column that is used by some constraint</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column that is used by some constraint</td>
</tr>
</tbody>
</table>
### 35.18. constraint_table_usage

The view `constraint_table_usage` identifies all tables in the current database that are used by some constraint and are owned by a currently enabled role. (This is different from the view `table_constraints`, which identifies all table constraints along with the table they are defined on.) For a foreign key constraint, this view identifies the table that the foreign key references. For a unique or primary key constraint, this view simply identifies the table the constraint belongs to. Check constraints and not-null constraints are not included in this view.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that is used by some constraint (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that is used by some constraint</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that is used by some constraint</td>
</tr>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
</tbody>
</table>

### 35.19. data_type_privileges

The view `data_type_privileges` identifies all data type descriptors that the current user has access to, by way of being the owner of the described object or having some privilege for it. A data type descriptor is generated whenever a data type is used in the definition of a table column, a domain, or a function (as parameter or return type) and stores some information about how the data type is used in that instance (for example, the declared maximum length, if applicable). Each data type descriptor is assigned an arbitrary identifier that is unique among the data type descriptor identifiers assigned for one object (table, domain, function). This view is probably not useful for applications, but it is used to define some other views in the information schema.
Table 35-17. data_type_privileges Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the described object (always the current database)</td>
</tr>
<tr>
<td>object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the described object</td>
</tr>
<tr>
<td>object_name</td>
<td>sql_identifier</td>
<td>Name of the described object</td>
</tr>
<tr>
<td>object_type</td>
<td>character_data</td>
<td>The type of the described object: one of TABLE (the data type descriptor pertains to a column of that table), DOMAIN (the data type descriptors pertains to that domain), ROUTINE (the data type descriptor pertains to a parameter or the return data type of that function).</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>The identifier of the data type descriptor, which is unique among the data type descriptors for that same object.</td>
</tr>
</tbody>
</table>

35.20. domain_constraints

The view domain_constraints contains all constraints belonging to domains defined in the current database. Only those domains are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-18. domain_constraints Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the domain</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
<tr>
<td>is_deferrable</td>
<td>yes_or_no</td>
<td>YES if the constraint is deferrable, NO if not</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>initially_deferred</td>
<td>yes_or_no</td>
<td>YES if the constraint is deferrable and initially deferred, NO if not</td>
</tr>
</tbody>
</table>

#### 35.21. domain_udt_usage

The view `domain_udt_usage` identifies all domains that are based on data types owned by a currently enabled role. Note that in PostgreSQL, built-in data types behave like user-defined types, so they are included here as well.

**Table 35-19. domain_udt_usage Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the domain data type is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the domain data type is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the domain data type</td>
</tr>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the domain</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
</tbody>
</table>

#### 35.22. domains

The view `domains` contains all domains defined in the current database. Only those domains are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 35-20. domains Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>domain_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the domain (always the current database)</td>
</tr>
<tr>
<td>domain_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the domain</td>
</tr>
<tr>
<td>domain_name</td>
<td>sql_identifier</td>
<td>Name of the domain</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the domain, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>If the domain has a character or bit string type, the declared maximum length; null for all other data types or if no maximum length was declared.</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>If the domain has a character type, the maximum possible length in octets (bytes) of a datum; null for all other data types. The maximum octet length depends on the declared character maximum length (see above) and the server encoding.</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the domain (always the current database), null if default or the data type of the domain is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the domain, null if default or the data type of the domain is not collatable</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the domain, null if default or the data type of the domain is not collatable</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>---------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>If the domain has a numeric type, this column contains the (declared or implicit) precision of the type for this domain. The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>If the domain has a numeric type, this column indicates in which base the values in the columns numeric_precision and numeric_scale are expressed. The value is either 2 or 10. For all other data types, this column is null.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>If the domain has an exact numeric type, this column contains the (declared or implicit) scale of the type for this domain. The scale indicates the number of significant digits to the right of the decimal point. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix. For all other data types, this column is null.</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>If data_type identifies a date, time, timestamp, or interval type, this column contains the (declared or implicit) fractional seconds precision of the type for this domain, that is, the number of decimal digits maintained following the decimal point in the seconds value. For all other data types, this column is null.</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>If <code>data_type</code> identifies an interval type, this column contains the specification which fields the intervals include for this domain, e.g., <code>YEAR TO MONTH</code>, <code>DAY TO SECOND</code>, etc. If no field restrictions were specified (that is, the interval accepts all fields), and for all other data types, this field is null.</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL (see <code>datetime_precision</code> for the fractional seconds precision of interval type domains)</td>
</tr>
<tr>
<td>domain_default</td>
<td>character_data</td>
<td>Default expression of the domain</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the domain data type is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the domain data type is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the domain data type</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the domain, unique among the data type descriptors pertaining to the domain (which is trivial, because a domain only contains one data type descriptor). This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
</tbody>
</table>
35.23. **element_types**

The view `element_types` contains the data type descriptors of the elements of arrays. When a table column, composite-type attribute, domain, function parameter, or function return value is defined to be of an array type, the respective information schema view only contains `ARRAY` in the column `data_type`. To obtain information on the element type of the array, you can join the respective view with this view. For example, to show the columns of a table with data types and array element types, if applicable, you could do:

```
SELECT c.column_name, c.data_type, e.data_type AS element_type
FROM information_schema.columns c LEFT JOIN information_schema.element_types e
    ON ((c.table_catalog, c.table_schema, c.table_name, 'TABLE', c.dtd_identifier)
        = (e.object_catalog, e.object_schema, e.object_name, e.object_type, e.collection_identifier))
WHERE c.table_schema = '...' AND c.table_name = '...
ORDER BY c.ordinal_position;
```

This view only includes objects that the current user has access to, by way of being the owner or having some privilege.

**Table 35-21. element_types Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the object that uses the array being described (always the current database)</td>
</tr>
<tr>
<td>object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the object that uses the array being described</td>
</tr>
<tr>
<td>object_name</td>
<td>sql_identifier</td>
<td>Name of the object that uses the array being described</td>
</tr>
<tr>
<td>object_type</td>
<td>character_data</td>
<td>The type of the object that uses the array being described: one of <code>TABLE</code> (the array is used by a column of that table), <code>USER-DEFINED TYPE</code> (the array is used by an attribute of that composite type), <code>DOMAIN</code> (the array is used by that domain), <code>ROUTINE</code> (the array is used by a parameter or the return data type of that function).</td>
</tr>
<tr>
<td>collection_type_identifier</td>
<td>sql_identifier</td>
<td>The identifier of the data type descriptor of the array being described. Use this to join with the <code>dtid_identifier</code> columns of other information schema views.</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------------------</td>
<td>-------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the array elements, if it is a built-in type, else USER-DEFINED (in that case, the type is identified in udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the collation of the element type (always the current database), null if default or the data type of the element is not collatable</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the collation of the element type, null if default or the data type of the element is not collatable</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Name of the collation of the element type, null if default or the data type of the element is not collatable</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to array element data types in PostgreSQL</td>
</tr>
<tr>
<td>domain_default</td>
<td>character_data</td>
<td>Not yet implemented</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the data type of the elements is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the data type of the elements is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the data type of the elements</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the element. This is currently not useful.</td>
</tr>
</tbody>
</table>

#### 35.24. enabled_roles

The view `enabled_roles` identifies the currently "enabled roles". The enabled roles are recursively defined as the current user together with all roles that have been granted to the enabled roles with automatic inheritance. In other words, these are all roles that the current user has direct or indirect, automatically inheriting membership in.

For permission checking, the set of “applicable roles” is applied, which can be broader than the set of enabled roles. So generally, it is better to use the view `applicable_roles` instead of this one; See Section 35.5 for details on `applicable_roles` view.
Table 35-22. enabled_roles Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>role_name</td>
<td>sql_identifier</td>
<td>Name of a role</td>
</tr>
</tbody>
</table>

35.25. foreign_data_wrapper_options

The view foreign_data_wrapper_options contains all the options defined for foreign-data wrappers in the current database. Only those foreign-data wrappers are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-23. foreign_data_wrapper_options Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_data_wrapper_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign-data wrapper is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_data_wrapper_name</td>
<td>sql_identifier</td>
<td>Name of the foreign-data wrapper</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

35.26. foreign_data_wrappers

The view foreign_data_wrappers contains all foreign-data wrappers defined in the current database. Only those foreign-data wrappers are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-24. foreign_data_wrappers Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_data_wrapper_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign-data wrapper (always the current database)</td>
</tr>
<tr>
<td>foreign_data_wrapper_name</td>
<td>sql_identifier</td>
<td>Name of the foreign-data wrapper</td>
</tr>
<tr>
<td>authorization_identifier</td>
<td>sql_identifier</td>
<td>Name of the owner of the foreign server</td>
</tr>
<tr>
<td>library_name</td>
<td>character_data</td>
<td>File name of the library that implementing this foreign-data wrapper</td>
</tr>
<tr>
<td>foreign_data_wrapper_language</td>
<td>character_data</td>
<td>Language used to implement this foreign-data wrapper</td>
</tr>
</tbody>
</table>
35.27. **foreign_server_options**

The view `foreign_server_options` contains all the options defined for foreign servers in the current database. Only those foreign servers are shown that the current user has access to (by way of being the owner or having some privilege).

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

35.28. **foreign_servers**

The view `foreign_servers` contains all foreign servers defined in the current database. Only those foreign servers are shown that the current user has access to (by way of being the owner or having some privilege).

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>foreign_data_wrapper_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign-data wrapper used by the foreign server (always the current database)</td>
</tr>
<tr>
<td>foreign_data_wrapper_name</td>
<td>sql_identifier</td>
<td>Name of the foreign-data wrapper used by the foreign server</td>
</tr>
<tr>
<td>foreign_server_type</td>
<td>character_data</td>
<td>Foreign server type information, if specified upon creation</td>
</tr>
<tr>
<td>foreign_server_version</td>
<td>character_data</td>
<td>Foreign server version information, if specified upon creation</td>
</tr>
<tr>
<td>authorization_identifier</td>
<td>sql_identifier</td>
<td>Name of the owner of the foreign server</td>
</tr>
</tbody>
</table>
Chapter 35. The Information Schema

35.29. foreign_table_options

The view foreign_table_options contains all the options defined for foreign tables in the current database. Only those foreign tables are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-27. foreign_table_options Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the foreign table (always the current database)</td>
</tr>
<tr>
<td>foreign_table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the foreign table</td>
</tr>
<tr>
<td>foreign_table_name</td>
<td>sql_identifier</td>
<td>Name of the foreign table</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option</td>
</tr>
</tbody>
</table>

35.30. foreign_tables

The view foreign_tables contains all foreign tables defined in the current database. Only those foreign tables are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-28. foreign_tables Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>foreign_table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign table is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the foreign table</td>
</tr>
<tr>
<td>foreign_table_name</td>
<td>sql_identifier</td>
<td>Name of the foreign table</td>
</tr>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server</td>
</tr>
</tbody>
</table>

35.31. key_column_usage

The view key_column_usage identifies all columns in the current database that are restricted by some unique, primary key, or foreign key constraint. Check constraints are not included in this view. Only those columns are shown that the current user has access to, by way of being the owner or having some privilege.
Chapter 35. The Information Schema

Table 35-29. key_column_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column that is restricted by this constraint (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column that is restricted by this constraint</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column that is restricted by this constraint</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column that is restricted by this constraint</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the column within the constraint key (count starts at 1)</td>
</tr>
<tr>
<td>position_in_unique_constraint</td>
<td>cardinal_number</td>
<td>For a foreign-key constraint, ordinal position of the referenced column within its unique constraint (count starts at 1); otherwise null</td>
</tr>
</tbody>
</table>

35.32. parameters

The view parameters contains information about the parameters (arguments) of all functions in the current database. Only those functions are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-30. parameters Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------</td>
<td>-----------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 35.40 for more information.</td>
</tr>
<tr>
<td>ordinal_position</td>
<td>cardinal_number</td>
<td>Ordinal position of the parameter in the argument list of the function (count starts at 1)</td>
</tr>
<tr>
<td>parameter_mode</td>
<td>character_data</td>
<td>IN for input parameter, OUT for output parameter, and INOUT for input/output parameter.</td>
</tr>
<tr>
<td>is_result</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>as_locator</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>parameter_name</td>
<td>sql_identifier</td>
<td>Name of the parameter, or null if the parameter has no name</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Data type of the parameter, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-----------------------------</td>
<td>----------------</td>
<td>----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to parameter data types in PostgreSQL</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the data type of the parameter is defined in (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the data type of the parameter is defined in</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the data type of the parameter</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the parameter, unique among the data type descriptors pertaining to the function. This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
<tr>
<td>parameter_default</td>
<td>character_data</td>
<td>The default expression of the parameter, or null if none or if the function is not owned by a currently enabled role.</td>
</tr>
</tbody>
</table>

### 35.33. referential_constraints

The view `referential_constraints` contains all referential (foreign key) constraints in the current database. Only those constraints are shown for which the current user has write access to the referencing table (by way of being the owner or having some privilege other than `SELECT`).

**Table 35-31. referential_constraints Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>unique_constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the unique or primary key constraint that the foreign key constraint references (always the current database)</td>
</tr>
<tr>
<td>unique_constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the unique or primary key constraint that the foreign key constraint references</td>
</tr>
<tr>
<td>unique_constraint_name</td>
<td>sql_identifier</td>
<td>Name of the unique or primary key constraint that the foreign key constraint references</td>
</tr>
<tr>
<td>match_option</td>
<td>character_data</td>
<td>Match option of the foreign key constraint: FULL, PARTIAL, or NONE.</td>
</tr>
</tbody>
</table>
Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>update_rule</td>
<td>character_data</td>
<td>Update rule of the foreign key constraint: CASCADE, SET NULL, SET DEFAULT,</td>
</tr>
<tr>
<td></td>
<td></td>
<td>RESTRICT, or NO ACTION.</td>
</tr>
<tr>
<td>delete_rule</td>
<td>character_data</td>
<td>Delete rule of the foreign key constraint: CASCADE, SET NULL, SET DEFAULT,</td>
</tr>
<tr>
<td></td>
<td></td>
<td>RESTRICT, or NO ACTION.</td>
</tr>
</tbody>
</table>

35.34. role_column_grants

The view `role_column_grants` identifies all privileges granted on columns where the grantor or grantee is a currently enabled role. Further information can be found under `column_privileges`. The only effective difference between this view and `column_privileges` is that this view omits columns that have been made accessible to the current user by way of a grant to `PUBLIC`.

Table 35-32. role_column_grants Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table that contains the column</td>
</tr>
<tr>
<td>column_name</td>
<td>sql_identifier</td>
<td>Name of the column</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Type of the privilege: SELECT, INSERT, UPDATE, or REFERENCES</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

35.35. role_routine_grants

The view `role_routine_grants` identifies all privileges granted on functions where the grantor or grantee is a currently enabled role. Further information can be found under `routine_privileges`. The only effective difference between this view and `routine_privileges` is that this view omits
functions that have been made accessible to the current user by way of a grant to PUBLIC.

### Table 35-33. `role_routine_grants` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td><code>sql_identifier</code></td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td><code>sql_identifier</code></td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>specific_catalog</td>
<td><code>sql_identifier</code></td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td><code>sql_identifier</code></td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td><code>sql_identifier</code></td>
<td>The “specific name” of the function. See Section 35.40 for more information.</td>
</tr>
<tr>
<td>routine_catalog</td>
<td><code>sql_identifier</code></td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>routine_schema</td>
<td><code>sql_identifier</code></td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>routine_name</td>
<td><code>sql_identifier</code></td>
<td>Name of the function (might be duplicated in case of overloading)</td>
</tr>
<tr>
<td>privilege_type</td>
<td><code>character_data</code></td>
<td>Always EXECUTE (the only privilege type for functions)</td>
</tr>
<tr>
<td>is_grantable</td>
<td><code>yes_or_no</code></td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

### 35.36. `role_table_grants`

The view `role_table_grants` identifies all privileges granted on tables or views where the grantor or grantee is a currently enabled role. Further information can be found under `table_privileges`. The only effective difference between this view and `table_privileges` is that this view omits tables that have been made accessible to the current user by way of a grant to PUBLIC.

### Table 35-34. `role_table_grants` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td><code>sql_identifier</code></td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td><code>sql_identifier</code></td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td><code>sql_identifier</code></td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
</tbody>
</table>
35.37. role_udt_grants

The view role_udt_grants is intended to identify USAGE privileges granted on user-defined types where the grantor or grantee is a currently enabled role. Further information can be found under udt_privileges. The only effective difference between this view and udt_privileges is that this view omits objects that have been made accessible to the current user by way of a grant to PUBLIC. Since data types do not have real privileges in PostgreSQL, but only an implicit grant to PUBLIC, this view is empty.

Table 35-35. role_udt_grants Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>The name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>The name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the type (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the type</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the type</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always TYPE USAGE</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>
35.38. role_usage_grants

The view role_usage_grants identifies USAGE privileges granted on various kinds of objects where the grantor or grantee is a currently enabled role. Further information can be found under usage_privileges. The only effective difference between this view and usage_privileges is that this view omits objects that have been made accessible to the current user by way of a grant to PUBLIC.

Table 35-36. role_usage_grants Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>The name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>The name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the object (always the current database)</td>
</tr>
<tr>
<td>object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the object, if applicable, else an empty string</td>
</tr>
<tr>
<td>object_name</td>
<td>sql_identifier</td>
<td>Name of the object</td>
</tr>
<tr>
<td>object_type</td>
<td>character_data</td>
<td>COLLATION or DOMAIN or FOREIGN DATA WRAPPER or FOREIGN SERVER or SEQUENCE</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always USAGE</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

35.39. routine_privileges

The view routine_privileges identifies all privileges granted on functions to a currently enabled role or by a currently enabled role. There is one row for each combination of function, grantor, and grantee.

Table 35-37. routine_privileges Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
</tbody>
</table>
### 35.40. routines

The view `routines` contains all functions in the current database. Only those functions are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-38. routines Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. This is a name that uniquely identifies the function in the schema, even if the real name of the function is overloaded. The format of the specific name is not defined, it should only be used to compare it to other instances of specific routine names.</td>
</tr>
<tr>
<td>routine_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>routine_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>routine_name</td>
<td>sql_identifier</td>
<td>Name of the function (might be duplicated in case of overloading)</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always <code>EXECUTE</code> (the only privilege type for functions)</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------</td>
<td>-----------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>routine_name</td>
<td>sql_identifier</td>
<td>Name of the function (might be duplicated in case of overloading)</td>
</tr>
<tr>
<td>routine_type</td>
<td>character_data</td>
<td>Always FUNCTION (In the future there might be other types of routines.)</td>
</tr>
<tr>
<td>module_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>module_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>module_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>Return data type of the function, if it is a built-in type, or ARRAY if it is some array (in that case, see the view element_types), else USER-DEFINED (in that case, the type is identified in type_udt_name and associated columns).</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------------</td>
<td>----------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>datetime_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_type</td>
<td>character_data</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>interval_precision</td>
<td>cardinal_number</td>
<td>Always null, since this information is not applied to return data types in PostgreSQL</td>
</tr>
<tr>
<td>type_udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the return data type of the function is defined in (always the current database)</td>
</tr>
<tr>
<td>type_udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that the return data type of the function is defined in</td>
</tr>
<tr>
<td>type_udt_name</td>
<td>sql_identifier</td>
<td>Name of the return data type of the function</td>
</tr>
<tr>
<td>scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>-----------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>maximum_cardinality</td>
<td>cardinal_number</td>
<td>Always null, because arrays always have unlimited maximum cardinality in PostgreSQL.</td>
</tr>
<tr>
<td>dtd_identifier</td>
<td>sql_identifier</td>
<td>An identifier of the data type descriptor of the return data type of this function, unique among the data type descriptors pertaining to the function. This is mainly useful for joining with other instances of such identifiers. (The specific format of the identifier is not defined and not guaranteed to remain the same in future versions.)</td>
</tr>
<tr>
<td>routine_body</td>
<td>character_data</td>
<td>If the function is an SQL function, then SQL, else EXTERNAL.</td>
</tr>
<tr>
<td>routine_definition</td>
<td>character_data</td>
<td>The source text of the function (null if the function is not owned by a currently enabled role). (According to the SQL standard, this column is only applicable if routine_body is SQL, but in PostgreSQL it will contain whatever source text was specified when the function was created.)</td>
</tr>
<tr>
<td>external_name</td>
<td>character_data</td>
<td>If this function is a C function, then the external name (link symbol) of the function; else null. (This works out to be the same value that is shown in routine_definition.)</td>
</tr>
<tr>
<td>external_language</td>
<td>character_data</td>
<td>The language the function is written in</td>
</tr>
<tr>
<td>parameter_style</td>
<td>character_data</td>
<td>Always GENERAL (The SQL standard defines other parameter styles, which are not available in PostgreSQL.)</td>
</tr>
<tr>
<td>is_deterministic</td>
<td>yes_or_no</td>
<td>If the function is declared immutable (called deterministic in the SQL standard), then YES, else NO. (You cannot query the other volatility levels available in PostgreSQL through the information schema.)</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sql_data_access</td>
<td>character_data</td>
<td>Always <strong>MODIFIES</strong>, meaning that the function possibly modifies SQL data. This information is not useful for PostgreSQL.</td>
</tr>
<tr>
<td>is_null_call</td>
<td>yes_or_no</td>
<td>If the function automatically returns null if any of its arguments are null, then <strong>YES</strong>, else <strong>NO</strong>.</td>
</tr>
<tr>
<td>sql_path</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>schema_levelroutine</td>
<td>yes_or_no</td>
<td>Always <strong>YES</strong> (The opposite would be a method of a user-defined type, which is a feature not available in PostgreSQL.)</td>
</tr>
<tr>
<td>max_dynamic_result_sets</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>is_user_defined_cast</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>isimplicitly_invocable</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>security_type</td>
<td>character_data</td>
<td>If the function runs with the privileges of the current user, then <strong>INVOKER</strong>, if the function runs with the privileges of the user who defined it, then <strong>DEFINER</strong>.</td>
</tr>
<tr>
<td>to_sql_specific_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>to_sql_specific_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>to_sql_specific_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>as_locator</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>created</td>
<td>time_stamp</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>last_altered</td>
<td>time_stamp</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>new_savepoint_level</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL.</td>
</tr>
<tr>
<td>is_udt_dependent</td>
<td>yes_or_no</td>
<td>Currently always <strong>NO</strong>. The alternative <strong>YES</strong> applies to a feature not available in PostgreSQL.</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>result_cast_from_data_type</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_as_locator</td>
<td>yes_or_no</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_max_length</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_octet_length</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_char_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_collation_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_collation_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_collation_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_numeric_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_numeric_precision_radix</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_numeric_scale</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_datetime_precision</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_interval_type</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_interval_precision</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_type_udt_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_type_udt_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_type_udt_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_scope_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_scope_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_scope_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>result_cast_maximum_cardinality</td>
<td>cardinal_number</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>result_cast_dtd_identifier</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>

#### 35.41. schemata

The view `schemata` contains all schemas in the current database that the current user has access to (by way of being the owner or having some privilege).

Table 35-39. `schemata` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>catalog_name</td>
<td>sql_identifier</td>
<td>Name of the database that the schema is contained in (always the current database)</td>
</tr>
<tr>
<td>schema_name</td>
<td>sql_identifier</td>
<td>Name of the schema</td>
</tr>
<tr>
<td>schema_owner</td>
<td>sql_identifier</td>
<td>Name of the owner of the schema</td>
</tr>
<tr>
<td>default_character_set_catalog</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>default_character_set_schema</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>default_character_set_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>sql_path</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>

#### 35.42. sequences

The view `sequences` contains all sequences defined in the current database. Only those sequences are shown that the current user has access to (by way of being the owner or having some privilege).

Table 35-40. `sequences` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sequence_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the sequence (always the current database)</td>
</tr>
<tr>
<td>sequence_schema</td>
<td>sql_identifier</td>
<td>Name of the schema</td>
</tr>
<tr>
<td>sequence_name</td>
<td>sql_identifier</td>
<td>Name of the sequence</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
<td>The data type of the sequence. In PostgreSQL, this is currently always bigint.</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
<td>This column contains the (declared or implicit) precision of the sequence data type (see above). The precision indicates the number of significant digits. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix.</td>
</tr>
<tr>
<td>numeric_precision_radix</td>
<td>cardinal_number</td>
<td>This column indicates in which base the values in the columns numeric_precision and numeric_scale are expressed. The value is either 2 or 10.</td>
</tr>
<tr>
<td>numeric_scale</td>
<td>cardinal_number</td>
<td>This column contains the (declared or implicit) scale of the sequence data type (see above). The scale indicates the number of significant digits to the right of the decimal point. It can be expressed in decimal (base 10) or binary (base 2) terms, as specified in the column numeric_precision_radix.</td>
</tr>
<tr>
<td>start_value</td>
<td>character_data</td>
<td>The start value of the sequence</td>
</tr>
<tr>
<td>minimum_value</td>
<td>character_data</td>
<td>The minimum value of the sequence</td>
</tr>
<tr>
<td>maximum_value</td>
<td>character_data</td>
<td>The maximum value of the sequence</td>
</tr>
<tr>
<td>increment</td>
<td>character_data</td>
<td>The increment of the sequence</td>
</tr>
<tr>
<td>cycle_option</td>
<td>yes_or_no</td>
<td>YES if the sequence cycles, else NO</td>
</tr>
</tbody>
</table>

Note that in accordance with the SQL standard, the start, minimum, maximum, and increment values are returned as character strings.

### 35.43. sql_features

The table sql_features contains information about which formal features defined in the SQL standard are supported by PostgreSQL. This is the same information that is presented in Appendix D. There you can also find some additional background information.
Table 35-41. sql_features Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>feature_id</td>
<td>character_data</td>
<td>Identifier string of the feature</td>
</tr>
<tr>
<td>feature_name</td>
<td>character_data</td>
<td>Descriptive name of the feature</td>
</tr>
<tr>
<td>sub_feature_id</td>
<td>character_data</td>
<td>Identifier string of the subfeature, or a zero-length string if not a subfeature</td>
</tr>
<tr>
<td>sub_feature_name</td>
<td>character_data</td>
<td>Descriptive name of the subfeature, or a zero-length string if not a subfeature</td>
</tr>
<tr>
<td>is_supported</td>
<td>yes_or_no</td>
<td>YES if the feature is fully supported by the current version of PostgreSQL, NO if not</td>
</tr>
<tr>
<td>is_verified_by</td>
<td>character_data</td>
<td>Always null, since the PostgreSQL development group does not perform formal testing of feature conformance</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment about the supported status of the feature</td>
</tr>
</tbody>
</table>

35.44. sql_implementation_info

The table sql_implementation_info contains information about various aspects that are left implementation-defined by the SQL standard. This information is primarily intended for use in the context of the ODBC interface; users of other interfaces will probably find this information to be of little use. For this reason, the individual implementation information items are not described here; you will find them in the description of the ODBC interface.

Table 35-42. sql_implementation_info Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>implementation_info_id</td>
<td>character_data</td>
<td>Identifier string of the implementation information item</td>
</tr>
<tr>
<td>implementation_info_name</td>
<td>character_data</td>
<td>Descriptive name of the implementation information item</td>
</tr>
<tr>
<td>integer_value</td>
<td>cardinal_number</td>
<td>Value of the implementation information item, or null if the value is contained in the column character_value</td>
</tr>
<tr>
<td>character_value</td>
<td>character_data</td>
<td>Value of the implementation information item, or null if the value is contained in the column integer_value</td>
</tr>
</tbody>
</table>
35.45. sql_languages

The table sql_languages contains one row for each SQL language binding that is supported by PostgreSQL. PostgreSQL supports direct SQL and embedded SQL in C; that is all you will learn from this table.

This table was removed from the SQL standard in SQL:2008, so there are no entries referring to standards later than SQL:2003.

Table 35-43. sql_languages Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sql_language_source</td>
<td>character_data</td>
<td>The name of the source of the language definition; always ISO 9075, that is, the SQL standard</td>
</tr>
<tr>
<td>sql_language_year</td>
<td>character_data</td>
<td>The year the standard referenced in sql_language_source was approved.</td>
</tr>
<tr>
<td>sql_language_conformance</td>
<td>character_data</td>
<td>The standard conformance level for the language binding. For ISO 9075:2003 this is always CORE.</td>
</tr>
<tr>
<td>sql_language_integrity</td>
<td>character_data</td>
<td>Always null (This value is relevant to an earlier version of the SQL standard.)</td>
</tr>
<tr>
<td>sql_language_implementation</td>
<td>character_data</td>
<td>Always null</td>
</tr>
<tr>
<td>sql_language_binding_style</td>
<td>character_data</td>
<td>The language binding style, either DIRECT or EMBEDDED</td>
</tr>
<tr>
<td>sql_language_programming_language</td>
<td>character_data</td>
<td>The programming language, if the binding style is EMBEDDED, else null. PostgreSQL only supports the language C.</td>
</tr>
</tbody>
</table>

35.46. sql_packages

The table sql_packages contains information about which feature packages defined in the SQL standard are supported by PostgreSQL. Refer to Appendix D for background information on feature packages.
Table 35-44. sql_packages Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>feature_id</td>
<td>character_data</td>
<td>Identifier string of the package</td>
</tr>
<tr>
<td>feature_name</td>
<td>character_data</td>
<td>Descriptive name of the package</td>
</tr>
<tr>
<td>is_supported</td>
<td>yes_or_no</td>
<td>YES if the package is fully supported by the current version of PostgreSQL, NO if not</td>
</tr>
<tr>
<td>is_verified_by</td>
<td>character_data</td>
<td>Always null, since the PostgreSQL development group does not perform formal testing of feature conformance</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment about the supported status of the package</td>
</tr>
</tbody>
</table>

35.47. sql_parts

The table sql_parts contains information about which of the several parts of the SQL standard are supported by PostgreSQL.

Table 35-45. sql_parts Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>feature_id</td>
<td>character_data</td>
<td>An identifier string containing the number of the part</td>
</tr>
<tr>
<td>feature_name</td>
<td>character_data</td>
<td>Descriptive name of the part</td>
</tr>
<tr>
<td>is_supported</td>
<td>yes_or_no</td>
<td>YES if the part is fully supported by the current version of PostgreSQL, NO if not</td>
</tr>
<tr>
<td>is_verified_by</td>
<td>character_data</td>
<td>Always null, since the PostgreSQL development group does not perform formal testing of feature conformance</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment about the supported status of the part</td>
</tr>
</tbody>
</table>

35.48. sql_sizing

The table sql_sizing contains information about various size limits and maximum values in PostgreSQL. This information is primarily intended for use in the context of the ODBC interface; users of other interfaces will probably find this information to be of little use. For this reason, the individual sizing items are not described here; you will find them in the description of the ODBC interface.

Table 35-46. sql_sizing Columns
### 35.49. sql_sizing_profiles

The table `sql_sizing_profiles` contains information about the `sql_sizing` values that are required by various profiles of the SQL standard. PostgreSQL does not track any SQL profiles, so this table is empty.

#### Table 35-47. sql_sizing_profiles Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sizing_id</td>
<td>cardinal_number</td>
<td>Identifier of the sizing item</td>
</tr>
<tr>
<td>sizing_name</td>
<td>character_data</td>
<td>Descriptive name of the sizing item</td>
</tr>
<tr>
<td>profile_id</td>
<td>character_data</td>
<td>Identifier string of a profile</td>
</tr>
<tr>
<td>required_value</td>
<td>cardinal_number</td>
<td>The value required by the SQL profile for the sizing item, or 0 if the profile places no limit on the sizing item, or null if the profile does not require any of the features for which the sizing item is applicable</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment pertaining to the sizing item within the profile</td>
</tr>
</tbody>
</table>

### 35.50. table_constraints

The view `table_constraints` contains all constraints belonging to tables that the current user owns or has some privilege other than `SELECT` on.

#### Table 35-48. table_constraints Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sizing_id</td>
<td>cardinal_number</td>
<td>Identifier of the sizing item</td>
</tr>
<tr>
<td>sizing_name</td>
<td>character_data</td>
<td>Descriptive name of the sizing item</td>
</tr>
<tr>
<td>profile_id</td>
<td>character_data</td>
<td>Identifier string of a profile</td>
</tr>
<tr>
<td>required_value</td>
<td>cardinal_number</td>
<td>The value required by the SQL profile for the sizing item, or 0 if the profile places no limit on the sizing item, or null if the profile does not require any of the features for which the sizing item is applicable</td>
</tr>
<tr>
<td>comments</td>
<td>character_data</td>
<td>Possibly a comment pertaining to the sizing item within the profile</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>constraint_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the constraint (always the current database)</td>
</tr>
<tr>
<td>constraint_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the constraint</td>
</tr>
<tr>
<td>constraint_name</td>
<td>sql_identifier</td>
<td>Name of the constraint</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>constraint_type</td>
<td>character_data</td>
<td>Type of the constraint: CHECK, FOREIGN KEY, PRIMARY KEY, or UNIQUE</td>
</tr>
<tr>
<td>is_deferrable</td>
<td>yes_or_no</td>
<td>YES if the constraint is deferrable, NO if not</td>
</tr>
<tr>
<td>initially_deferred</td>
<td>yes_or_no</td>
<td>YES if the constraint is deferrable and initially deferred, NO if not</td>
</tr>
</tbody>
</table>

#### 35.51. `table_privileges`

The view `table_privileges` identifies all privileges granted on tables or views to a currently enabled role or by a currently enabled role. There is one row for each combination of table, grantor, and grantee.

**Table 35-49, `table_privileges` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Type of the privilege: SELECT, INSERT, UPDATE, DELETE, TRUNCATE, REFERENCES, or TRIGGER</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
<tr>
<td>with_hierarchy</td>
<td>yes_or_no</td>
<td>In the SQL standard, WITH HIERARCHY OPTION is a separate (sub-)privilege allowing certain operations on table inheritance hierarchies. In PostgreSQL, this is included in the SELECT privilege, so this column shows YES if the privilege is SELECT, else NO.</td>
</tr>
</tbody>
</table>

#### 35.52. tables

The view `tables` contains all tables and views defined in the current database. Only those tables and views are shown that the current user has access to (by way of being the owner or having some privilege).

**Table 35-50. tables Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the table</td>
</tr>
<tr>
<td>table_type</td>
<td>character_data</td>
<td>Type of the table: BASE TABLE for a persistent base table (the normal table type), VIEW for a view, FOREIGN TABLE for a foreign table, or LOCAL TEMPORARY for a temporary table</td>
</tr>
<tr>
<td>self_referencing_column_name</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>reference_generation</td>
<td>character_data</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>user_defined_type_catalog</td>
<td>sql_identifier</td>
<td>If the table is a typed table, the name of the database that contains the underlying data type (always the current database), else null.</td>
</tr>
</tbody>
</table>
### 35.53. transforms

The view `transforms` contains information about the transforms defined in the current database. More precisely, it contains a row for each function contained in a transform (the “from SQL” or “to SQL” function).

#### Table 35-51. transforms Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the type the transform is for (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the type the transform is for</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the type the transform is for</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 35.40 for more information.</td>
</tr>
</tbody>
</table>
### 35.54. triggered_update_columns

For triggers in the current database that specify a column list (like `UPDATE OF column1, column2`), the view `triggered_update_columns` identifies these columns. Triggers that do not specify a column list are not included in this view. Only those columns are shown that the current user owns or has some privilege other than `SELECT` on.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>trigger_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the trigger (always the current database)</td>
</tr>
<tr>
<td>trigger_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the trigger</td>
</tr>
<tr>
<td>trigger_name</td>
<td>sql_identifier</td>
<td>Name of the trigger</td>
</tr>
<tr>
<td>event_object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that the trigger is defined on (always the current database)</td>
</tr>
<tr>
<td>event_object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that the trigger is defined on</td>
</tr>
<tr>
<td>event_object_table</td>
<td>sql_identifier</td>
<td>Name of the table that the trigger is defined on</td>
</tr>
<tr>
<td>event_object_column</td>
<td>sql_identifier</td>
<td>Name of the column that the trigger is defined on</td>
</tr>
</tbody>
</table>

### 35.55. triggers

The view `triggers` contains all triggers defined in the current database on tables and views that the current user owns or has some privilege other than `SELECT` on.
Table 35-53. triggers Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>trigger_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the trigger (always the current database)</td>
</tr>
<tr>
<td>trigger_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the trigger</td>
</tr>
<tr>
<td>trigger_name</td>
<td>sql_identifier</td>
<td>Name of the trigger</td>
</tr>
<tr>
<td>event_manipulation</td>
<td>character_data</td>
<td>Event that fires the trigger (INSERT, UPDATE, or DELETE)</td>
</tr>
<tr>
<td>event_object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that the trigger is defined on (always the current database)</td>
</tr>
<tr>
<td>event_object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that the trigger is defined on</td>
</tr>
<tr>
<td>event_object_table</td>
<td>sql_identifier</td>
<td>Name of the table that the trigger is defined on</td>
</tr>
<tr>
<td>action_order</td>
<td>cardinal_number</td>
<td>Not yet implemented</td>
</tr>
<tr>
<td>action_condition</td>
<td>character_data</td>
<td>WHEN condition of the trigger, null if none (also null if the table is not owned by a currently enabled role)</td>
</tr>
<tr>
<td>action_statement</td>
<td>character_data</td>
<td>Statement that is executed by the trigger (currently always EXECUTE PROCEDURE function(...))</td>
</tr>
<tr>
<td>action_orientation</td>
<td>character_data</td>
<td>Identifies whether the trigger fires once for each processed row or once for each statement (ROW or STATEMENT)</td>
</tr>
<tr>
<td>action_timing</td>
<td>character_data</td>
<td>Time at which the trigger fires (BEFORE, AFTER, or INSTEAD OF)</td>
</tr>
<tr>
<td>action_reference_old_table</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>action_reference_new_table</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>action_reference_old_row</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>action_reference_new_row</td>
<td>sql_identifier</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
<tr>
<td>created</td>
<td>time_stamp</td>
<td>Applies to a feature not available in PostgreSQL</td>
</tr>
</tbody>
</table>

Triggers in PostgreSQL have two incompatibilities with the SQL standard that affect the
representation in the information schema. First, trigger names are local to each table in PostgreSQL, rather than being independent schema objects. Therefore there can be duplicate trigger names defined in one schema, so long as they belong to different tables. (trigger_catalog and trigger_schema are really the values pertaining to the table that the trigger is defined on.) Second, triggers can be defined to fire on multiple events in PostgreSQL (e.g., ON INSERT OR UPDATE), whereas the SQL standard only allows one. If a trigger is defined to fire on multiple events, it is represented as multiple rows in the information schema, one for each type of event. As a consequence of these two issues, the primary key of the view triggers is really (trigger_catalog, trigger_schema, event_object_table, trigger_name, event_manipulation) instead of (trigger_catalog, trigger_schema, trigger_name), which is what the SQL standard specifies. Nonetheless, if you define your triggers in a manner that conforms with the SQL standard (trigger names unique in the schema and only one event type per trigger), this will not affect you.

**Note:** Prior to PostgreSQL 9.1, this view's columns action_timing, action_reference_old_table, action_reference_new_table, action_reference_old_row, and action_reference_new_row were named condition_timing, condition_reference_old_table, condition_reference_new_table, condition_reference_old_row, and condition_reference_new_row respectively. That was how they were named in the SQL:1999 standard. The new naming conforms to SQL:2003 and later.

### 35.56. udt_privileges

The view udt_privileges identifies USAGE privileges granted on user-defined types to a currently enabled role or by a currently enabled role. There is one row for each combination of type, grantor, and grantee. This view shows only composite types (see under Section 35.58 for why); see Section 35.57 for domain privileges.

**Table 35-54. udt_privileges Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>udt_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the type (always the current database)</td>
</tr>
<tr>
<td>udt_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the type</td>
</tr>
<tr>
<td>udt_name</td>
<td>sql_identifier</td>
<td>Name of the type</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always TYPE USAGE</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>
35.57. usage_privileges

The view `usage_privileges` identifies USAGE privileges granted on various kinds of objects to a currently enabled role or by a currently enabled role. In PostgreSQL, this currently applies to collations, domains, foreign-data wrappers, foreign servers, and sequences. There is one row for each combination of object, grantor, and grantee.

Since collations do not have real privileges in PostgreSQL, this view shows implicit non-grantable USAGE privileges granted by the owner to PUBLIC for all collations. The other object types, however, show real privileges.

In PostgreSQL, sequences also support SELECT and UPDATE privileges in addition to the USAGE privilege. These are nonstandard and therefore not visible in the information schema.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>grantor</td>
<td>sql_identifier</td>
<td>Name of the role that granted the privilege</td>
</tr>
<tr>
<td>grantee</td>
<td>sql_identifier</td>
<td>Name of the role that the privilege was granted to</td>
</tr>
<tr>
<td>object_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the object (always the current database)</td>
</tr>
<tr>
<td>object_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the object, if applicable, else an empty string</td>
</tr>
<tr>
<td>object_name</td>
<td>sql_identifier</td>
<td>Name of the object</td>
</tr>
<tr>
<td>object_type</td>
<td>character_data</td>
<td>Collation or domain or foreign data wrapper or foreign server or sequence</td>
</tr>
<tr>
<td>privilege_type</td>
<td>character_data</td>
<td>Always USAGE</td>
</tr>
<tr>
<td>is_grantable</td>
<td>yes_or_no</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
</tbody>
</table>

35.58. user_defined_types

The view `user_defined_types` currently contains all composite types defined in the current database. Only those types are shown that the current user has access to (by way of being the owner or having some privilege).

SQL knows about two kinds of user-defined types: structured types (also known as composite types in PostgreSQL) and distinct types (not implemented in PostgreSQL). To be future-proof, use the column `user_defined_type_category` to differentiate between these. Other user-defined types such as base types and enums, which are PostgreSQL extensions, are not shown here. For domains, see Section 35.22 instead.

<table>
<thead>
<tr>
<th>Name</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>privilege_type</td>
<td>Always USAGE</td>
</tr>
<tr>
<td>is_grantable</td>
<td>YES if the privilege is grantable, NO if not</td>
</tr>
<tr>
<td>Name</td>
<td>Data Type</td>
</tr>
<tr>
<td>----------------------------------</td>
<td>---------------</td>
</tr>
<tr>
<td>user_defined_type_catalog</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>user_defined_type_schema</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>user_defined_type_name</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>user_defined_type_category</td>
<td>character_data</td>
</tr>
<tr>
<td>is_instantiable</td>
<td>yes_or_no</td>
</tr>
<tr>
<td>is_final</td>
<td>yes_or_no</td>
</tr>
<tr>
<td>ordering_form</td>
<td>character_data</td>
</tr>
<tr>
<td>ordering_category</td>
<td>character_data</td>
</tr>
<tr>
<td>orderingRoutine_catalog</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>orderingRoutine_schema</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>orderingRoutine_name</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>reference_type</td>
<td>character_data</td>
</tr>
<tr>
<td>data_type</td>
<td>character_data</td>
</tr>
<tr>
<td>character_maximum_length</td>
<td>cardinal_number</td>
</tr>
<tr>
<td>character_octet_length</td>
<td>cardinal_number</td>
</tr>
<tr>
<td>character_set_catalog</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>character_set_schema</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>character_set_name</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>collation_catalog</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>collation_schema</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>collation_name</td>
<td>sql_identifier</td>
</tr>
<tr>
<td>numeric_precision</td>
<td>cardinal_number</td>
</tr>
</tbody>
</table>
### 35.59. user_mapping_options

The view `user_mapping_options` contains all the options defined for user mappings in the current database. Only those user mappings are shown where the current user has access to the corresponding foreign server (by way of being the owner or having some privilege).

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>authorization_identifier</td>
<td>sql_identifier</td>
<td>Name of the user being mapped, or <code>PUBLIC</code> if the mapping is public</td>
</tr>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server used by this mapping is defined in (always the current database)</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server used by this mapping</td>
</tr>
<tr>
<td>option_name</td>
<td>sql_identifier</td>
<td>Name of an option</td>
</tr>
<tr>
<td>option_value</td>
<td>character_data</td>
<td>Value of the option. This column will show as null unless the current user is the user being mapped, or the mapping is for <code>PUBLIC</code> and the current user is the server owner, or the current user is a superuser. The intent is to protect password information stored as user mapping option.</td>
</tr>
</tbody>
</table>

The table above lists the columns that comprise the view `user_mapping_options`. Each column represents a specific aspect of the user mappings in the database.
35.60. **user_mappings**

The view `user_mappings` contains all user mappings defined in the current database. Only those user mappings are shown where the current user has access to the corresponding foreign server (by way of being the owner or having some privilege).

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>authorization_identifier</td>
<td>sql_identifier</td>
<td>Name of the user being mapped, or PUBLIC if the mapping is public</td>
</tr>
<tr>
<td>foreign_server_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that the foreign server used by this mapping is defined in</td>
</tr>
<tr>
<td>foreign_server_name</td>
<td>sql_identifier</td>
<td>Name of the foreign server used by this mapping</td>
</tr>
</tbody>
</table>

35.61. **view_column_usage**

The view `view_column_usage` identifies all columns that are used in the query expression of a view (the `SELECT` statement that defines the view). A column is only included if the table that contains the column is owned by a currently enabled role.

**Note:** Columns of system tables are not included. This should be fixed sometime.

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>view_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the view (always the current database)</td>
</tr>
<tr>
<td>view_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the view</td>
</tr>
<tr>
<td>view_name</td>
<td>sql_identifier</td>
<td>Name of the view</td>
</tr>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database that contains the table that contains the column that is used by the view (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema that contains the table that contains the column that is used by the view</td>
</tr>
</tbody>
</table>

978
### Table 35-60. view_routine_usage Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the view (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the view</td>
</tr>
<tr>
<td>table_name</td>
<td>sql_identifier</td>
<td>Name of the view</td>
</tr>
<tr>
<td>specific_catalog</td>
<td>sql_identifier</td>
<td>Name of the database containing the function (always the current database)</td>
</tr>
<tr>
<td>specific_schema</td>
<td>sql_identifier</td>
<td>Name of the schema containing the function</td>
</tr>
<tr>
<td>specific_name</td>
<td>sql_identifier</td>
<td>The “specific name” of the function. See Section 35.40 for more information.</td>
</tr>
</tbody>
</table>

### 35.63. view_table_usage

The view view_table_usage identifies all tables that are used in the query expression of a view (the SELECT statement that defines the view). A table is only included if that table is owned by a currently enabled role.

**Note:** System tables are not included. This should be fixed sometime.
### 35.64. views

The view `views` contains all views defined in the current database. Only those views are shown that the current user has access to (by way of being the owner or having some privilege).

#### Table 35-62. views Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_catalog</td>
<td><code>sql_identifier</code></td>
<td>Name of the database that contains the view (always the current database)</td>
</tr>
<tr>
<td>table_schema</td>
<td><code>sql_identifier</code></td>
<td>Name of the schema that contains the view</td>
</tr>
<tr>
<td>table_name</td>
<td><code>sql_identifier</code></td>
<td>Name of the view</td>
</tr>
<tr>
<td>view_definition</td>
<td><code>character_data</code></td>
<td>Query expression defining the view (null if the view is not owned by a currently enabled role)</td>
</tr>
<tr>
<td>check_option</td>
<td><code>character_data</code></td>
<td>CASCaded or LOCAL if the view has a CHECK OPTION defined on it, NONE if not</td>
</tr>
<tr>
<td>is_updatable</td>
<td><code>yes_or_no</code></td>
<td>YES if the view is updatable (allows UPDATE and DELETE), NO if not</td>
</tr>
<tr>
<td>is_insertable_into</td>
<td><code>yes_or_no</code></td>
<td>YES if the view is insertable into (allows INSERT), NO if not</td>
</tr>
</tbody>
</table>
### Chapter 35. The Information Schema

<table>
<thead>
<tr>
<th>Name</th>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>is_trigger_updatable</td>
<td>yes_or_no</td>
<td>YES if the view has an INSTEAD OF UPDATE trigger defined on it, NO if not</td>
</tr>
<tr>
<td>is_trigger_deletable</td>
<td>yes_or_no</td>
<td>YES if the view has an INSTEAD OF DELETE trigger defined on it, NO if not</td>
</tr>
<tr>
<td>is_trigger_insertable_int</td>
<td>yes_or_no</td>
<td>YES if the view has an INSTEAD OF INSERT trigger defined on it, NO if not</td>
</tr>
</tbody>
</table>
V. Server Programming

This part is about extending the server functionality with user-defined functions, data types, triggers, etc. These are advanced topics which should probably be approached only after all the other user documentation about PostgreSQL has been understood. Later chapters in this part describe the server-side programming languages available in the PostgreSQL distribution as well as general issues concerning server-side programming languages. It is essential to read at least the earlier sections of Chapter 36 (covering functions) before diving into the material about server-side programming languages.
Chapter 36. Extending SQL

In the sections that follow, we will discuss how you can extend the PostgreSQL SQL query language by adding:

- functions (starting in Section 36.3)
- aggregates (starting in Section 36.10)
- data types (starting in Section 36.11)
- operators (starting in Section 36.12)
- operator classes for indexes (starting in Section 36.14)
- packages of related objects (starting in Section 36.15)

36.1. How Extensibility Works

PostgreSQL is extensible because its operation is catalog-driven. If you are familiar with standard relational database systems, you know that they store information about databases, tables, columns, etc., in what are commonly known as system catalogs. (Some systems call this the data dictionary.) The catalogs appear to the user as tables like any other, but the DBMS stores its internal bookkeeping in them. One key difference between PostgreSQL and standard relational database systems is that PostgreSQL stores much more information in its catalogs: not only information about tables and columns, but also information about data types, functions, access methods, and so on. These tables can be modified by the user, and since PostgreSQL bases its operation on these tables, this means that PostgreSQL can be extended by users. By comparison, conventional database systems can only be extended by changing hardcoded procedures in the source code or by loading modules specially written by the DBMS vendor.

The PostgreSQL server can moreover incorporate user-written code into itself through dynamic loading. That is, the user can specify an object code file (e.g., a shared library) that implements a new type or function, and PostgreSQL will load it as required. Code written in SQL is even more trivial to add to the server. This ability to modify its operation “on the fly” makes PostgreSQL uniquely suited for rapid prototyping of new applications and storage structures.

36.2. The PostgreSQL Type System

PostgreSQL data types are divided into base types, composite types, domains, and pseudo-types.

36.2.1. Base Types

Base types are those, like `int4`, that are implemented below the level of the SQL language (typically in a low-level language such as C). They generally correspond to what are often known as abstract data types. PostgreSQL can only operate on such types through functions provided by the user and only understands the behavior of such types to the extent that the user describes them. Base types are further subdivided into scalar and array types. For each scalar type, a corresponding array type is automatically created that can hold variable-size arrays of that scalar type.
36.2.2. Composite Types

Composite types, or row types, are created whenever the user creates a table. It is also possible to use CREATE TYPE to define a “stand-alone” composite type with no associated table. A composite type is simply a list of types with associated field names. A value of a composite type is a row or record of field values. The user can access the component fields from SQL queries. Refer to Section 8.16 for more information on composite types.

36.2.3. Domains

A domain is based on a particular base type and for many purposes is interchangeable with its base type. However, a domain can have constraints that restrict its valid values to a subset of what the underlying base type would allow.

Domains can be created using the SQL command CREATE DOMAIN. Their creation and use is not discussed in this chapter.

36.2.4. Pseudo-Types

There are a few “pseudo-types” for special purposes. Pseudo-types cannot appear as columns of tables or attributes of composite types, but they can be used to declare the argument and result types of functions. This provides a mechanism within the type system to identify special classes of functions. Table 8-25 lists the existing pseudo-types.

36.2.5. Polymorphic Types

Five pseudo-types of special interest are anyelement, anyarray, anynonarray, anyenum, and anyrange, which are collectively called polymorphic types. Any function declared using these types is said to be a polymorphic function. A polymorphic function can operate on many different data types, with the specific data type(s) being determined by the data types actually passed to it in a particular call.

Polymorphic arguments and results are tied to each other and are resolved to a specific data type when a query calling a polymorphic function is parsed. Each position (either argument or return value) declared as anyelement is allowed to have any specific actual data type, but in any given call they must all be the same actual type. Each position declared as anyarray can have any array data type, but similarly they must all be the same type. And similarly, positions declared as anyrange must all be the same range type. Furthermore, if there are positions declared anyarray and others declared anyelement, the actual array type in the anyarray positions must be an array whose elements are the same type appearing in the anyelement positions. Similarly, if there are positions declared anyrange and others declared anyelement, the actual range type in the anyrange positions must be a range whose subtype is the same type appearing in the anyelement positions. anynonarray is treated exactly the same as anyelement, but adds the additional constraint that the actual type must be a range type. anyrange is treated exactly the same as anyelement, but adds the additional constraint that the actual type must be an array type.

Thus, when more than one argument position is declared with a polymorphic type, the net effect is that only certain combinations of actual argument types are allowed. For example, a function declared as equal(anyelement, anyelement) will take any two input values, so long as they are of the same data type.
When the return value of a function is declared as a polymorphic type, there must be at least one argument position that is also polymorphic, and the actual data type supplied as the argument determines the actual result type for that call. For example, if there were not already an array subscripting mechanism, one could define a function that implements subscripting as `subscript(anyarray, integer) returns anyelement`. This declaration constrains the actual first argument to be an array type, and allows the parser to infer the correct result type from the actual first argument’s type. Another example is that a function declared as `f(anyarray) returns anyenum` will only accept arrays of enum types.

Note that `anynonarray` and `anyenum` do not represent separate type variables; they are the same type as `anyelement`, just with an additional constraint. For example, declaring a function as `f(anyelement, anyenum)` is equivalent to declaring it as `f(anyenum, anyenum)`: both actual arguments have to be the same enum type.

A variadic function (one taking a variable number of arguments, as in Section 36.4.5) can be polymorphic: this is accomplished by declaring its last parameter as `VARIADIC anyarray`. For purposes of argument matching and determining the actual result type, such a function behaves the same as if you had written the appropriate number of `anynonarray` parameters.

### 36.3. User-defined Functions

PostgreSQL provides four kinds of functions:

- query language functions (functions written in SQL) (Section 36.4)
- procedural language functions (functions written in, for example, PL/pgSQL or PL/Tcl) (Section 36.7)
- internal functions (Section 36.8)
- C-language functions (Section 36.9)

Every kind of function can take base types, composite types, or combinations of these as arguments (parameters). In addition, every kind of function can return a base type or a composite type. Functions can also be defined to return sets of base or composite values.

Many kinds of functions can take or return certain pseudo-types (such as polymorphic types), but the available facilities vary. Consult the description of each kind of function for more details.

It’s easiest to define SQL functions, so we’ll start by discussing those. Most of the concepts presented for SQL functions will carry over to the other types of functions.

Throughout this chapter, it can be useful to look at the reference page of the CREATE FUNCTION command to understand the examples better. Some examples from this chapter can be found in `funcs.sql` and `funcs.c` in the `src/tutorial` directory in the PostgreSQL source distribution.

### 36.4. Query Language (SQL) Functions

SQL functions execute an arbitrary list of SQL statements, returning the result of the last query in the list. In the simple (non-set) case, the first row of the last query’s result will be returned. (Bear in
mind that “the first row” of a multirow result is not well-defined unless you use ORDER BY.) If the last query happens to return no rows at all, the null value will be returned.

Alternatively, an SQL function can be declared to return a set (that is, multiple rows) by specifying the function’s return type as SETOF sometype, or equivalently by declaring it as RETURNS TABLE (columns). In this case all rows of the last query’s result are returned. Further details appear below.

The body of an SQL function must be a list of SQL statements separated by semicolons. A semicolon after the last statement is optional. Unless the function is declared to return void, the last statement must be a SELECT, or an INSERT, UPDATE, or DELETE that has a RETURNING clause.

Any collection of commands in the SQL language can be packaged together and defined as a function. Besides SELECT queries, the commands can include data modification queries (INSERT, UPDATE, and DELETE), as well as other SQL commands. (You cannot use transaction control commands, e.g., COMMIT, SAVEPOINT, and some utility commands, e.g. VACUUM, in SQL functions.) However, the final command must be a SELECT or have a RETURNING clause that returns whatever is specified as the function’s return type. Alternatively, if you want to define a SQL function that performs actions but has no useful value to return, you can define it as returning void. For example, this function removes rows with negative salaries from the emp table:

```
CREATE FUNCTION clean_emp() RETURNS void AS '
  DELETE FROM emp
   WHERE salary < 0;
' LANGUAGE SQL;
```

```
SELECT clean_emp();
```

```
clean_emp
------------
(1 row)
```

**Note:** The entire body of a SQL function is parsed before any of it is executed. While a SQL function can contain commands that alter the system catalogs (e.g., CREATE TABLE), the effects of such commands will not be visible during parse analysis of later commands in the function. Thus, for example, `CREATE TABLE foo (...); INSERT INTO foo VALUES(...);` will not work as desired if packaged up into a single SQL function, since foo won’t exist yet when the INSERT command is parsed. It’s recommended to use PL/PgSQL instead of a SQL function in this type of situation.

The syntax of the CREATE FUNCTION command requires the function body to be written as a string constant. It is usually most convenient to use dollar quoting (see Section 4.1.2.4) for the string constant. If you choose to use regular single-quoted string constant syntax, you must double single quote marks (‘) and backslashes (\) (assuming escape string syntax) in the body of the function (see Section 4.1.2.1).

### 36.4.1. Arguments for SQL Functions

Arguments of a SQL function can be referenced in the function body using either names or numbers. Examples of both methods appear below.
To use a name, declare the function argument as having a name, and then just write that name in the function body. If the argument name is the same as any column name in the current SQL command within the function, the column name will take precedence. To override this, qualify the argument name with the name of the function itself, that is function_name.argument_name. (If this would conflict with a qualified column name, again the column name wins. You can avoid the ambiguity by choosing a different alias for the table within the SQL command.)

In the older numeric approach, arguments are referenced using the syntax $n$: $1$ refers to the first input argument, $2$ to the second, and so on. This will work whether or not the particular argument was declared with a name.

If an argument is of a composite type, then the dot notation, e.g., argname.fieldname or $1.fieldname$, can be used to access attributes of the argument. Again, you might need to qualify the argument’s name with the function name to make the form with an argument name unambiguous.

SQL function arguments can only be used as data values, not as identifiers. Thus for example this is reasonable:

```
INSERT INTO mytable VALUES ($1);
```

but this will not work:

```
INSERT INTO $1 VALUES (42);
```

**Note:** The ability to use names to reference SQL function arguments was added in PostgreSQL 9.2. Functions to be used in older servers must use the $n$ notation.

### 36.4.2. SQL Functions on Base Types

The simplest possible SQL function has no arguments and simply returns a base type, such as integer:

```
CREATE FUNCTION one() RETURNS integer AS $$
  SELECT 1 AS result;
$$ LANGUAGE SQL;
```

```
-- Alternative syntax for string literal:
CREATE FUNCTION one() RETURNS integer AS 'SELECT 1 AS result';
LANGUAGE SQL;
```

```
SELECT one();
```

```
one
-----
1
```

Notice that we defined a column alias within the function body for the result of the function (with the name result), but this column alias is not visible outside the function. Hence, the result is labeled one instead of result.
It is almost as easy to define SQL functions that take base types as arguments:

```
CREATE FUNCTION add_em(x integer, y integer) RETURNS integer AS $$
SELECT x + y;
$$ LANGUAGE SQL;
```

```
SELECT add_em(1, 2) AS answer;
```

```
answer
--------
3
``` 

Alternatively, we could dispense with names for the arguments and use numbers:

```
CREATE FUNCTION add_em(integer, integer) RETURNS integer AS $$
SELECT $1 + $2;
$$ LANGUAGE SQL;
```

```
SELECT add_em(1, 2) AS answer;
```

```
answer
--------
3
``` 

Here is a more useful function, which might be used to debit a bank account:

```
CREATE FUNCTION tf1 (accountno integer, debit numeric) RETURNS integer AS $$
UPDATE bank
    SET balance = balance - debit
WHERE accountno = tf1.accountno;
SELECT 1;
$$ LANGUAGE SQL;
```

A user could execute this function to debit account 17 by $100.00 as follows:

```
SELECT tf1(17, 100.0);
```

In this example, we chose the name `accountno` for the first argument, but this is the same as the name of a column in the `bank` table. Within the `UPDATE` command, `accountno` refers to the column `bank.accountno`, so `tf1.accountno` must be used to refer to the argument. We could of course avoid this by using a different name for the argument.

In practice one would probably like a more useful result from the function than a constant 1, so a more likely definition is:

```
CREATE FUNCTION tf1 (accountno integer, debit numeric) RETURNS integer AS $$
UPDATE bank
    SET balance = balance - debit
WHERE accountno = tf1.accountno;
SELECT balance FROM bank WHERE accountno = tf1.accountno;
$$ LANGUAGE SQL;
```
which adjusts the balance and returns the new balance. The same thing could be done in one command using \texttt{RETURNING}:

\begin{verbatim}
CREATE FUNCTION tf1 (accountno integer, debit numeric) RETURNS integer AS $$
UPDATE bank
  SET balance = balance - debit
WHERE accountno = tf1.accountno
RETURNING balance;
$$ LANGUAGE SQL;
\end{verbatim}

36.4.3. SQL Functions on Composite Types

When writing functions with arguments of composite types, we must not only specify which argument we want but also the desired attribute (field) of that argument. For example, suppose that \texttt{emp} is a table containing employee data, and therefore also the name of the composite type of each row of the table. Here is a function \texttt{double_salary} that computes what someone’s salary would be if it were doubled:

\begin{verbatim}
CREATE TABLE emp (  
  name text,  
  salary numeric,  
  age integer,  
  cubicle point  
);  
INSERT INTO emp VALUES ('Bill', 4200, 45, '(2,1)');
CREATE FUNCTION double_salary(emp) RETURNS numeric AS $$
SELECT $1.salary * 2 AS salary;
$$ LANGUAGE SQL;
SELECT name, double_salary(emp.*) AS dream 
FROM emp 
WHERE emp.cubicle ~= point '(2,1)';
\end{verbatim}

```
name | dream
------+-------
Bill | 8400
```

Notice the use of the syntax $1.salary to select one field of the argument row value. Also notice how the calling \texttt{SELECT} command uses \texttt{table_name.*} to select the entire current row of a table as a composite value. The table row can alternatively be referenced using just the table name, like this:

\begin{verbatim}
SELECT name, double_salary(emp) AS dream 
FROM emp 
WHERE emp.cubicle ~= point '(2,1)';
\end{verbatim}

but this usage is deprecated since it’s easy to get confused. (See Section 8.16.5 for details about these two notations for the composite value of a table row.)

Sometimes it is handy to construct a composite argument value on-the-fly. This can be done with the \texttt{ROW} construct. For example, we could adjust the data being passed to the function:
SELECT name, double_salary(ROW(name, salary*1.1, age, cubicle)) AS dream
FROM emp;

It is also possible to build a function that returns a composite type. This is an example of a function
that returns a single emp row:

CREATE FUNCTION new_emp() RETURNS emp AS $$
SELECT text 'None' AS name,
    1000.0 AS salary,
    25 AS age,
    point '(2,2)' AS cubicle;
$$ LANGUAGE SQL;

In this example we have specified each of the attributes with a constant value, but any computation
could have been substituted for these constants.

Note two important things about defining the function:

• The select list order in the query must be exactly the same as that in which the columns appear in
  the table associated with the composite type. (Naming the columns, as we did above, is irrelevant
to the system.)

• You must typecast the expressions to match the definition of the composite type, or you will get
  errors like this:

    ERROR: function declared to return emp returns varchar instead of text at column 1

A different way to define the same function is:

CREATE FUNCTION new_emp() RETURNS emp AS $$
    SELECT ROW('None', 1000.0, 25, '(2,2)')::emp;
$$ LANGUAGE SQL;

Here we wrote a SELECT that returns just a single column of the correct composite type. This isn’t
really better in this situation, but it is a handy alternative in some cases — for example, if we need to
compute the result by calling another function that returns the desired composite value.

We could call this function directly either by using it in a value expression:

SELECT new_emp();

    new_emp
    ---------------
    (None,1000.0,25,"(2,2)")

or by calling it as a table function:

SELECT * FROM new_emp();

    name | salary | age | cubicle
    -------+--------+-----+---------
    None   | 1000.0 | 25  | (2,2)

The second way is described more fully in Section 36.4.7.

When you use a function that returns a composite type, you might want only one field (attribute) from
its result. You can do that with syntax like this:
SELECT (new_emp()).name;

name
------
None

The extra parentheses are needed to keep the parser from getting confused. If you try to do it without them, you get something like this:

SELECT new_emp().name;
ERROR: syntax error at or near "."
LINE 1: SELECT new_emp().name;

Another option is to use functional notation for extracting an attribute:

SELECT name(new_emp());

name
------
None

As explained in Section 8.16.5, the field notation and functional notation are equivalent.

Another way to use a function returning a composite type is to pass the result to another function that accepts the correct row type as input:

CREATE FUNCTION getname(emp) RETURNS text AS $$
    SELECT $1.name;
$$ LANGUAGE SQL;

SELECT getname(new_emp());

getname
----------
None
(1 row)

36.4.4. SQL Functions with Output Parameters

An alternative way of describing a function’s results is to define it with output parameters, as in this example:

CREATE FUNCTION add_em (IN x int, IN y int, OUT sum int) AS 'SELECT x + y'
LANGUAGE SQL;

SELECT add_em(3,7);
add_em
-------
  10
(1 row)
This is not essentially different from the version of `add_em` shown in Section 36.4.2. The real value of output parameters is that they provide a convenient way of defining functions that return several columns. For example,

```
CREATE FUNCTION sum_n_product (x int, y int, OUT sum int, OUT product int)
AS 'SELECT x + y, x * y'
LANGUAGE SQL;
```

```
SELECT * FROM sum_n_product(11,42);
```

```
sum | product
-----+---------
  53 |   462
```

What has essentially happened here is that we have created an anonymous composite type for the result of the function. The above example has the same end result as

```
CREATE TYPE sum_prod AS (sum int, product int);
CREATE FUNCTION sum_n_product (int, int) RETURNS sum_prod
AS 'SELECT $1 + $2, $1 * $2'
LANGUAGE SQL;
```

but not having to bother with the separate composite type definition is often handy. Notice that the names attached to the output parameters are not just decoration, but determine the column names of the anonymous composite type. (If you omit a name for an output parameter, the system will choose a name on its own.)

Notice that output parameters are not included in the calling argument list when invoking such a function from SQL. This is because PostgreSQL considers only the input parameters to define the function’s calling signature. That means also that only the input parameters matter when referencing the function for purposes such as dropping it. We could drop the above function with either of

```
DROP FUNCTION sum_n_product (x int, y int, OUT sum int, OUT product int);
DROP FUNCTION sum_n_product (int, int);
```

Parameters can be marked as `IN` (the default), `OUT`, `INOUT`, or `VARIADIC`. An `INOUT` parameter serves as both an input parameter (part of the calling argument list) and an output parameter (part of the result record type). `VARIADIC` parameters are input parameters, but are treated specially as described next.

### 36.4.5. SQL Functions with Variable Numbers of Arguments

SQL functions can be declared to accept variable numbers of arguments, so long as all the “optional” arguments are of the same data type. The optional arguments will be passed to the function as an array. The function is declared by marking the last parameter as `VARIADIC`; this parameter must be declared as being of an array type. For example:

```
CREATE FUNCTION mleast(VARIADIC arr numeric[]) RETURNS numeric AS $$
SELECT min($1[i]) FROM generate_subscripts($1, 1) g(i);
$$ LANGUAGE SQL;
```

```
SELECT mleast(10, -1, 5, 4.4);
```

993
Effectively, all the actual arguments at or beyond the \texttt{VARIADIC} position are gathered up into a one-dimensional array, as if you had written

\begin{verbatim}
SELECT mleast(ARRAY[10, -1, 5, 4.4]); -- doesn’t work
\end{verbatim}

You can’t actually write that, though — or at least, it will not match this function definition. A parameter marked \texttt{VARIADIC} matches one or more occurrences of its element type, not of its own type. Sometimes it is useful to be able to pass an already-constructed array to a variadic function; this is particularly handy when one variadic function wants to pass on its array parameter to another one. Also, this is the only secure way to call a variadic function found in a schema that permits untrusted users to create objects; see Section 10.3. You can do this by specifying \texttt{VARIADIC} in the call:

\begin{verbatim}
SELECT mleast(VARIADIC ARRAY[10, -1, 5, 4.4]);
\end{verbatim}

This prevents expansion of the function’s variadic parameter into its element type, thereby allowing the array argument value to match normally. \texttt{VARIADIC} can only be attached to the last actual argument of a function call. Specifying \texttt{VARIADIC} in the call is also the only way to pass an empty array to a variadic function, for example:

\begin{verbatim}
SELECT mleast(VARIADIC ARRAY[]::numeric[]);
\end{verbatim}

Simply writing \texttt{SELECT mleast()} does not work because a variadic parameter must match at least one actual argument. (You could define a second function also named \texttt{mleast}, with no parameters, if you wanted to allow such calls.)

The array element parameters generated from a variadic parameter are treated as not having any names of their own. This means it is not possible to call a variadic function using named arguments (Section 4.3), except when you specify \texttt{VARIADIC}. For example, this will work:

\begin{verbatim}
SELECT mleast(VARIADIC arr =&gt; ARRAY[10, -1, 5, 4.4]);
\end{verbatim}

but not these:

\begin{verbatim}
SELECT mleast(arr =&gt; 10);
SELECT mleast(arr =&gt; ARRAY[10, -1, 5, 4.4]);
\end{verbatim}

### 36.4.6. SQL Functions with Default Values for Arguments

Functions can be declared with default values for some or all input arguments. The default values are inserted whenever the function is called with insufficiently many actual arguments. Since arguments can only be omitted from the end of the actual argument list, all parameters after a parameter with a default value have to have default values as well. (Although the use of named argument notation could allow this restriction to be relaxed, it’s still enforced so that positional argument notation works sensibly.) Whether or not you use it, this capability creates a need for precautions when calling functions in databases where some users mistrust other users; see Section 10.3.
For example:

```sql
CREATE FUNCTION foo(a int, b int DEFAULT 2, c int DEFAULT 3)
RETURNS int
LANGUAGE SQL
AS $$
SELECT $1 + $2 + $3;
$$;

SELECT foo(10, 20, 30);
foo
-----
60
(1 row)

SELECT foo(10, 20);
foo
-----
33
(1 row)

SELECT foo(10);
foo
-----
15
(1 row)

SELECT foo(); -- fails since there is no default for the first argument
ERROR: function foo() does not exist
```

The `=` sign can also be used in place of the key word `DEFAULT`.

### 36.4.7. SQL Functions as Table Sources

All SQL functions can be used in the `FROM` clause of a query, but it is particularly useful for functions returning composite types. If the function is defined to return a base type, the table function produces a one-column table. If the function is defined to return a composite type, the table function produces a column for each attribute of the composite type.

Here is an example:

```sql
CREATE TABLE foo (fooid int, foosubid int, fooname text);
INSERT INTO foo VALUES (1, 1, 'Joe');
INSERT INTO foo VALUES (1, 2, 'Ed');
INSERT INTO foo VALUES (2, 1, 'Mary');

CREATE FUNCTION getfoo(int) RETURNS foo AS $$
SELECT * FROM foo WHERE fooid = $1;
$$ LANGUAGE SQL;

SELECT *., upper(fooname) FROM getfoo(1) AS t1;
```

<table>
<thead>
<tr>
<th>fooid</th>
<th>foosubid</th>
<th>fooname</th>
<th>upper</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Joe</td>
<td>JOE</td>
</tr>
</tbody>
</table>

(1 row)
As the example shows, we can work with the columns of the function’s result just the same as if they were columns of a regular table.

Note that we only got one row out of the function. This is because we did not use SETOF. That is described in the next section.

### 36.4.8. SQL Functions Returning Sets

When an SQL function is declared as returning SETOF someType, the function’s final query is executed to completion, and each row it outputs is returned as an element of the result set.

This feature is normally used when calling the function in the FROM clause. In this case each row returned by the function becomes a row of the table seen by the query. For example, assume that table foo has the same contents as above, and we say:

```sql
CREATE FUNCTION getfoo(int) RETURNS SETOF foo AS $$
    SELECT * FROM foo WHERE fooid = $1;
$$ LANGUAGE SQL;
SELECT * FROM getfoo(1) AS t1;
```

Then we would get:

<table>
<thead>
<tr>
<th>fooid</th>
<th>foosubid</th>
<th>fooname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Joe</td>
</tr>
<tr>
<td>1</td>
<td>2</td>
<td>Ed</td>
</tr>
</tbody>
</table>
(2 rows)

It is also possible to return multiple rows with the columns defined by output parameters, like this:

```sql
CREATE TABLE tab (y int, z int);
INSERT INTO tab VALUES (1, 2), (3, 4), (5, 6), (7, 8);
CREATE FUNCTION sum_n_product_with_tab (x int, OUT sum int, OUT product int) RETURNS SETOF record AS $$
    SELECT $1 + tab.y, $1 * tab.y FROM tab;
$$ LANGUAGE SQL;
SELECT * FROM sum_n_product_with_tab(10);
```

<table>
<thead>
<tr>
<th>sum</th>
<th>product</th>
</tr>
</thead>
<tbody>
<tr>
<td>11</td>
<td>10</td>
</tr>
<tr>
<td>13</td>
<td>30</td>
</tr>
<tr>
<td>15</td>
<td>50</td>
</tr>
<tr>
<td>17</td>
<td>70</td>
</tr>
</tbody>
</table>
(4 rows)

The key point here is that you must write RETURNS SETOF record to indicate that the function returns multiple rows instead of just one. If there is only one output parameter, write that parameter’s type instead of record.

It is frequently useful to construct a query’s result by invoking a set-returning function multiple times, with the parameters for each invocation coming from successive rows of a table or subquery. The
preferred way to do this is to use the LATERAL key word, which is described in Section 7.2.1.5. Here is an example using a set-returning function to enumerate elements of a tree structure:

```
SELECT * FROM nodes;
name    | parent
---------+--------
Top      |        
Child1   | Top    
Child2   | Top    
Child3   | Top    
SubChild1| Child1 
SubChild2| Child1 
(6 rows)
```

```
CREATE FUNCTION listchildren(text) RETURNS SETOF text AS $$
SELECT name FROM nodes WHERE parent = $1
$$ LANGUAGE SQL STABLE;
```

```
SELECT * FROM listchildren('Top');
listchildren
------------
Child1
Child2
Child3
(3 rows)
```

```
SELECT name, child FROM nodes, LATERAL listchildren(name) AS child;
name    | child
--------+-------
Top     | Child1
Top     | Child2
Top     | Child3
Child1  | SubChild1
Child1  | SubChild2
(5 rows)
```

This example does not do anything that we couldn’t have done with a simple join, but in more complex calculations the option to put some of the work into a function can be quite convenient.

Currently, functions returning sets can also be called in the select list of a query. For each row that the query generates by itself, the function returning set is invoked, and an output row is generated for each element of the function’s result set. Note, however, that this capability is deprecated and might be removed in future releases. The previous example could also be done with queries like these:

```
SELECT listchildren('Top');
listchildren
------------
Child1
Child2
Child3
(3 rows)
```

```
SELECT name, listchildren(name) FROM nodes;
name    | listchildren
--------+--------------
Top     | Child1
Top     | Child2
997
```

997
In the last SELECT, notice that no output row appears for Child2, Child3, etc. This happens because listchildren returns an empty set for those arguments, so no result rows are generated. This is the same behavior as we got from an inner join to the function result when using the LATERAL syntax.

**Note:** If a function’s last command is `INSERT`, `UPDATE`, or `DELETE` with `RETURNING`, that command will always be executed to completion, even if the function is not declared with `SETOF` or the calling query does not fetch all the result rows. Any extra rows produced by the `RETURNING` clause are silently dropped, but the commanded table modifications still happen (and are all completed before returning from the function).

**Note:** The key problem with using set-returning functions in the select list, rather than the `FROM` clause, is that putting more than one set-returning function in the same select list does not behave very sensibly. (What you actually get if you do so is a number of output rows equal to the least common multiple of the numbers of rows produced by each set-returning function.) The LATERAL syntax produces less surprising results when calling multiple set-returning functions, and should usually be used instead.

### 36.4.9. SQL Functions Returning `TABLE`

There is another way to declare a function as returning a set, which is to use the syntax `RETURNS TABLE(columns)`. This is equivalent to using one or more `OUT` parameters plus marking the function as returning `SETOF record` (or `SETOF` a single output parameter’s type, as appropriate). This notation is specified in recent versions of the SQL standard, and thus may be more portable than using `SETOF`.

For example, the preceding sum-and-product example could also be done this way:

```sql
CREATE FUNCTION sum_n_product_with_tab (x int)
RETURNS TABLE(sum int, product int) AS $$
SELECT $1 + tab.y, $1 * tab.y FROM tab;
$$ LANGUAGE SQL;
```

It is not allowed to use explicit `OUT` or `INOUT` parameters with the `RETURNS TABLE` notation — you must put all the output columns in the `TABLE` list.

### 36.4.10. Polymorphic SQL Functions

SQL functions can be declared to accept and return the polymorphic types `anyelement`, `anyarray`, `anynonarray`, `anyenum`, and `anyrange`. See Section 36.2.5 for a more detailed explanation of polymorphic functions. Here is a polymorphic function `make_array` that builds up an array from two arbitrary data type elements:

```sql
CREATE FUNCTION make_array(anyelement, anyelement) RETURNS anyarray AS $$
SELECT ARRAY[$1, $2];
$$ LANGUAGE SQL;
```
SELECT make_array(1, 2) AS intarray, make_array('a':text, 'b') AS textarray;

<p>| | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>intarray</td>
<td>textarray</td>
</tr>
<tr>
<td>{1,2}</td>
<td>{a,b}</td>
</tr>
</tbody>
</table>

(1 row)

Notice the use of the typecast 'a':text to specify that the argument is of type text. This is required if the argument is just a string literal, since otherwise it would be treated as type unknown, and array of unknown is not a valid type. Without the typecast, you will get errors like this:

```
ERROR: could not determine polymorphic type because input has type "unknown"
```

It is permitted to have polymorphic arguments with a fixed return type, but the converse is not. For example:

CREATE FUNCTION is_greater(anyelement, anyelement) RETURNS boolean AS $$
SELECT $1 > $2;
$$ LANGUAGE SQL;

SELECT is_greater(1, 2);

<table>
<thead>
<tr>
<th>is_greater</th>
</tr>
</thead>
<tbody>
<tr>
<td>f</td>
</tr>
</tbody>
</table>

CREATE FUNCTION invalid_func() RETURNS anyelement AS $$
SELECT 1;
$$ LANGUAGE SQL;

ERROR: cannot determine result data type
DETAIL: A function returning a polymorphic type must have at least one polymorphic argument.

Polymorphism can be used with functions that have output arguments. For example:

CREATE FUNCTION dup (f1 anyelement, OUT f2 anyelement, OUT f3 anyarray) AS 'select $1, array[$1,$1]' LANGUAGE SQL;

SELECT * FROM dup(22);

<table>
<thead>
<tr>
<th>f2</th>
<th>f3</th>
</tr>
</thead>
<tbody>
<tr>
<td>22</td>
<td>{22,22}</td>
</tr>
</tbody>
</table>

(1 row)

Polymorphism can also be used with variadic functions. For example:

CREATE FUNCTION anyleast (VARIADIC anyarray) RETURNS anyelement AS $$
SELECT min($1[i]) FROM generate_subscripts($1, 1) g(i);
$$ LANGUAGE SQL;

SELECT anyleast(10, -1, 5, 4);

<table>
<thead>
<tr>
<th>anyleast</th>
</tr>
</thead>
<tbody>
<tr>
<td>999</td>
</tr>
</tbody>
</table>
Chapter 36. Extending SQL

---

(1 row)

SELECT anyleast('abc':text, 'def');
anyleast
---
abc
(1 row)

CREATE FUNCTION concat_values(text, VARIADIC anyarray) RETURNS text AS $$
SELECT array_to_string($2, $1);
$$ LANGUAGE SQL;

SELECT concat_values('|', 1, 4, 2);
cocat_values
---
1|4|2
(1 row)

36.4.11. SQL Functions with Collations

When a SQL function has one or more parameters of collatable data types, a collation is identified for each function call depending on the collations assigned to the actual arguments, as described in Section 23.2. If a collation is successfully identified (i.e., there are no conflicts of implicit collations among the arguments) then all the collatable parameters are treated as having that collation implicitly. This will affect the behavior of collation-sensitive operations within the function. For example, using the anyleast function described above, the result of

SELECT anyleast('abc':text, 'ABC');

will depend on the database’s default collation. In C locale the result will be ABC, but in many other locales it will be abc. The collation to use can be forced by adding a COLLATE clause to any of the arguments, for example

SELECT anyleast('abc':text, 'ABC' COLLATE "C");

Alternatively, if you wish a function to operate with a particular collation regardless of what it is called with, insert COLLATE clauses as needed in the function definition. This version of anyleast would always use en_US locale to compare strings:

CREATE FUNCTION anyleast (VARIADIC anyarray) RETURNS anyelement AS $$
SELECT min($1[i] COLLATE "en_US") FROM generate_subscripts($1, 1) g(i);
$$ LANGUAGE SQL;

But note that this will throw an error if applied to a non-collatable data type.

If no common collation can be identified among the actual arguments, then a SQL function treats its parameters as having their data types’ default collation (which is usually the database’s default collation, but could be different for parameters of domain types).

The behavior of collatable parameters can be thought of as a limited form of polymorphism, applicable only to textual data types.
36.5. Function Overloading

More than one function can be defined with the same SQL name, so long as the arguments they take are different. In other words, function names can be overloaded. Whether or not you use it, this capability entails security precautions when calling functions in databases where some users mistrust other users; see Section 10.3. When a query is executed, the server will determine which function to call from the data types and the number of the provided arguments. Overloading can also be used to simulate functions with a variable number of arguments, up to a finite maximum number.

When creating a family of overloaded functions, one should be careful not to create ambiguities. For instance, given the functions:

```sql
CREATE FUNCTION test(int, real) RETURNS ...  
CREATE FUNCTION test(smallint, double precision) RETURNS ...  
```

it is not immediately clear which function would be called with some trivial input like `test(1, 1.5)`. The currently implemented resolution rules are described in Chapter 10, but it is unwise to design a system that subtly relies on this behavior.

A function that takes a single argument of a composite type should generally not have the same name as any attribute (field) of that type. Recall that `attribute(table)` is considered equivalent to `table.attribute`. In the case that there is an ambiguity between a function on a composite type and an attribute of the composite type, the attribute will always be used. It is possible to override that choice by schema-qualifying the function name (that is, `schema.attribute`) but it’s better to avoid the problem by not choosing conflicting names.

Another possible conflict is between variadic and non-variadic functions. For instance, it is possible to create both `foo(numeric)` and `foo(VARIADIC numeric[])`. In this case it is unclear which one should be matched to a call providing a single numeric argument, such as `foo(10.1)`. The rule is that the function appearing earlier in the search path is used, or if the two functions are in the same schema, the non-variadic one is preferred.

When overloading C-language functions, there is an additional constraint: The C name of each function in the family of overloaded functions must be different from the C names of all other functions, either internal or dynamically loaded. If this rule is violated, the behavior is not portable. You might get a run-time linker error, or one of the functions will get called (usually the internal one). The alternative form of the `AS` clause for the SQL `CREATE FUNCTION` command decouples the SQL function name from the function name in the C source code. For instance:

```sql
CREATE FUNCTION test(int) RETURNS int  
    AS 'filename', 'test_1arg'  
    LANGUAGE C;  
CREATE FUNCTION test(int, int) RETURNS int  
    AS 'filename', 'test_2arg'  
    LANGUAGE C;  
```

The names of the C functions here reflect one of many possible conventions.

36.6. Function Volatility Categories

Every function has a volatility classification, with the possibilities being VOLATILE, STABLE, or IMmutable. VOLATILE is the default if the `CREATE FUNCTION` command does not specify a
category. The volatility category is a promise to the optimizer about the behavior of the function:

- A **VOLATILE** function can do anything, including modifying the database. It can return different results on successive calls with the same arguments. The optimizer makes no assumptions about the behavior of such functions. A query using a volatile function will re-evaluate the function at every row where its value is needed.

- A **STABLE** function cannot modify the database and is guaranteed to return the same results given the same arguments for all rows within a single statement. This category allows the optimizer to optimize multiple calls of the function to a single call. In particular, it is safe to use an expression containing such a function in an index scan condition. (Since an index scan will evaluate the comparison value only once, not once at each row, it is not valid to use a **VOLATILE** function in an index scan condition.)

- An **IMMUTABLE** function cannot modify the database and is guaranteed to return the same results given the same arguments forever. This category allows the optimizer to pre-evaluate the function when a query calls it with constant arguments. For example, a query like `SELECT ... WHERE x = 2 + 2` can be simplified on sight to `SELECT ... WHERE x = 4`, because the function underlying the integer addition operator is marked **IMMUTABLE**.

For best optimization results, you should label your functions with the strictest volatility category that is valid for them.

Any function with side-effects must be labeled **VOLATILE**, so that calls to it cannot be optimized away. Even a function with no side-effects needs to be labeled **VOLATILE** if its value can change within a single query; some examples are `random()`, `currval()`, `timeofday()`.

Another important example is that the `current_timestamp` family of functions qualify as **STABLE**, since their values do not change within a transaction.

There is relatively little difference between **STABLE** and **IMMUTABLE** categories when considering simple interactive queries that are planned and immediately executed: it doesn’t matter a lot whether a function is executed once during planning or once during query execution startup. But there is a big difference if the plan is saved and reused later. Labeling a function **IMMUTABLE** when it really isn’t might allow it to be prematurely folded to a constant during planning, resulting in a stale value being re-used during subsequent uses of the plan. This is a hazard when using prepared statements or when using function languages that cache plans (such as PL/pgSQL).

For functions written in SQL or in any of the standard procedural languages, there is a second important property determined by the volatility category, namely the visibility of any data changes that have been made by the SQL command that is calling the function. A **VOLATILE** function will see such changes, a **STABLE** or **IMMUTABLE** function will not. This behavior is implemented using the snapshotting behavior of MVCC (see Chapter 13): **STABLE** and **IMMUTABLE** functions use a snapshot established as of the start of the calling query, whereas **VOLATILE** functions obtain a fresh snapshot at the start of each query they execute.

**Note:** Functions written in C can manage snapshots however they want, but it’s usually a good idea to make C functions work this way too.

Because of this snapshotting behavior, a function containing only `SELECT` commands can safely be marked **STABLE**, even if it selects from tables that might be undergoing modifications by concurrent
queries. PostgreSQL will execute all commands of a \texttt{STABLE} function using the snapshot established for the calling query, and so it will see a fixed view of the database throughout that query.

The same snapshotting behavior is used for \texttt{SELECT} commands within \texttt{IMMUTABLE} functions. It is generally unwise to select from database tables within an \texttt{IMMUTABLE} function at all, since the immutability will be broken if the table contents ever change. However, PostgreSQL does not enforce that you do not do that.

A common error is to label a function \texttt{IMMUTABLE} when its results depend on a configuration parameter. For example, a function that manipulates timestamps might well have results that depend on the TimeZone setting. For safety, such functions should be labeled \texttt{STABLE} instead.

\textbf{Note:} PostgreSQL requires that \texttt{STABLE} and \texttt{IMMUTABLE} functions contain no SQL commands other than \texttt{SELECT} to prevent data modification. (This is not a completely bulletproof test, since such functions could still call \texttt{VOLATILE} functions that modify the database. If you do that, you will find that the \texttt{STABLE} or \texttt{IMMUTABLE} function does not notice the database changes applied by the called function, since they are hidden from its snapshot.)

\section{36.7. Procedural Language Functions}

PostgreSQL allows user-defined functions to be written in other languages besides SQL and C. These other languages are generically called \texttt{procedural languages} (PLs). Procedural languages aren’t built into the PostgreSQL server; they are offered by loadable modules. See Chapter 40 and following chapters for more information.

\section{36.8. Internal Functions}

Internal functions are functions written in C that have been statically linked into the PostgreSQL server. The “body” of the function definition specifies the C-language name of the function, which need not be the same as the name being declared for SQL use. (For reasons of backward compatibility, an empty body is accepted as meaning that the C-language function name is the same as the SQL name.)

Normally, all internal functions present in the server are declared during the initialization of the database cluster (see Section 18.2), but a user could use \texttt{CREATE FUNCTION} to create additional alias names for an internal function. Internal functions are declared in \texttt{CREATE FUNCTION} with language name \texttt{internal}. For instance, to create an alias for the $\sqrt{x}$ function:

\begin{verbatim}
CREATE FUNCTION square_root(double precision) RETURNS double precision
  AS 'dsqrt'
  LANGUAGE internal
  STRICT;
\end{verbatim}

(Most internal functions expect to be declared “strict”.)

\textbf{Note:} Not all “predefined” functions are “internal” in the above sense. Some predefined functions are written in SQL.
36.9. C-Language Functions

User-defined functions can be written in C (or a language that can be made compatible with C, such as C++). Such functions are compiled into dynamically loadable objects (also called shared libraries) and are loaded by the server on demand. The dynamic loading feature is what distinguishes “C language” functions from “internal” functions — the actual coding conventions are essentially the same for both. (Hence, the standard internal function library is a rich source of coding examples for user-defined C functions.)

Two different calling conventions are currently used for C functions. The newer “version 1” calling convention is indicated by writing a `PG_FUNCTION_INFO_V1()` macro call for the function, as illustrated below. Lack of such a macro indicates an old-style (“version 0”) function. The language name specified in `CREATE FUNCTION` is `C` in either case. Old-style functions are now deprecated because of portability problems and lack of functionality, but they are still supported for compatibility reasons.

36.9.1. Dynamic Loading

The first time a user-defined function in a particular loadable object file is called in a session, the dynamic loader loads that object file into memory so that the function can be called. The `CREATE FUNCTION` for a user-defined C function must therefore specify two pieces of information for the function: the name of the loadable object file, and the C name (link symbol) of the specific function to call within that object file. If the C name is not explicitly specified then it is assumed to be the same as the SQL function name.

The following algorithm is used to locate the shared object file based on the name given in the `CREATE FUNCTION` command:

1. If the name is an absolute path, the given file is loaded.
2. If the name starts with the string `$libdir`, that part is replaced by the PostgreSQL package library directory name, which is determined at build time.
3. If the name does not contain a directory part, the file is searched for in the path specified by the configuration variable `dynamic_library_path`.
4. Otherwise (the file was not found in the path, or it contains a non-absolute directory part), the dynamic loader will try to take the name as given, which will most likely fail. (It is unreliable to depend on the current working directory.)

If this sequence does not work, the platform-specific shared library file name extension (often `.so`) is appended to the given name and this sequence is tried again. If that fails as well, the load will fail.

It is recommended to locate shared libraries either relative to `$libdir` or through the dynamic library path. This simplifies version upgrades if the new installation is at a different location. The actual directory that `$libdir` stands for can be found out with the command `pg_config --pkglibdir`.

The user ID the PostgreSQL server runs as must be able to traverse the path to the file you intend to load. Making the file or a higher-level directory not readable and/or not executable by the postgres user is a common mistake.

In any case, the file name that is given in the `CREATE FUNCTION` command is recorded literally in the system catalogs, so if the file needs to be loaded again the same procedure is applied.

**Note:** PostgreSQL will not compile a C function automatically. The object file must be compiled before it is referenced in a `CREATE FUNCTION` command. See Section 36.9.6 for additional information.
To ensure that a dynamically loaded object file is not loaded into an incompatible server, PostgreSQL checks that the file contains a “magic block” with the appropriate contents. This allows the server to detect obvious incompatibilities, such as code compiled for a different major version of PostgreSQL. A magic block is required as of PostgreSQL 8.2. To include a magic block, write this in one (and only one) of the module source files, after having included the header fmgr.h:

```c
#ifdef PG_MODULE_MAGIC
PG_MODULE_MAGIC;
#endif
```

The `#ifdef` test can be omitted if the code doesn’t need to compile against pre-8.2 PostgreSQL releases.

After it is used for the first time, a dynamically loaded object file is retained in memory. Future calls in the same session to the function(s) in that file will only incur the small overhead of a symbol table lookup. If you need to force a reload of an object file, for example after recompiling it, begin a fresh session.

Optionally, a dynamically loaded file can contain initialization and finalization functions. If the file includes a function named `_PG_init`, that function will be called immediately after loading the file. The function receives no parameters and should return void. If the file includes a function named `_PG_fini`, that function will be called immediately before unloading the file. Likewise, the function receives no parameters and should return void. Note that `_PG_fini` will only be called during an unload of the file, not during process termination. (Presently, unloads are disabled and will never occur, but this may change in the future.)

### 36.9.2. Base Types in C-Language Functions

To know how to write C-language functions, you need to know how PostgreSQL internally represents base data types and how they can be passed to and from functions. Internally, PostgreSQL regards a base type as a “blob of memory”. The user-defined functions that you define over a type in turn define the way that PostgreSQL can operate on it. That is, PostgreSQL will only store and retrieve the data from disk and use your user-defined functions to input, process, and output the data.

Base types can have one of three internal formats:

- pass by value, fixed-length
- pass by reference, fixed-length
- pass by reference, variable-length

By-value types can only be 1, 2, or 4 bytes in length (also 8 bytes, if `sizeof(Datum)` is 8 on your machine). You should be careful to define your types such that they will be the same size (in bytes) on all architectures. For example, the `long` type is dangerous because it is 4 bytes on some machines and 8 bytes on others, whereas `int` type is 4 bytes on most Unix machines. A reasonable implementation of the `int4` type on Unix machines might be:

```c
/* 4-byte integer, passed by value */
typedef int int4;
```
(The actual PostgreSQL C code calls this type \texttt{int32}, because it is a convention in C that \texttt{int}\texttt{XX} means \texttt{XX} \texttt{bits}. Note therefore also that the C type \texttt{int8} is 1 byte in size. The SQL type \texttt{int8} is called \texttt{int64} in C. See also Table 36-1.)

On the other hand, fixed-length types of any size can be passed by-reference. For example, here is a sample implementation of a PostgreSQL type:

```c
/* 16-byte structure, passed by reference */
typedef struct {
    double x, y;
} Point;
```

Only pointers to such types can be used when passing them in and out of PostgreSQL functions. To return a value of such a type, allocate the right amount of memory with \texttt{palloc}, fill in the allocated memory, and return a pointer to it. (Also, if you just want to return the same value as one of your input arguments that’s of the same data type, you can skip the extra \texttt{palloc} and just return the pointer to the input value.)

Finally, all variable-length types must also be passed by reference. All variable-length types must begin with an opaque length field of exactly 4 bytes, which will be set by \texttt{SET_VARSIZE}; never set this field directly! All data to be stored within that type must be located in the memory immediately following that length field. The length field contains the total length of the structure, that is, it includes the size of the length field itself.

Another important point is to avoid leaving any uninitialized bits within data type values; for example, take care to zero out any alignment padding bytes that might be present in structs. Without this, logically-equivalent constants of your data type might be seen as unequal by the planner, leading to inefficient (though not incorrect) plans.

**Warning**

*Never* modify the contents of a pass-by-reference input value. If you do so you are likely to corrupt on-disk data, since the pointer you are given might point directly into a disk buffer. The sole exception to this rule is explained in Section 36.10.

As an example, we can define the type \texttt{text} as follows:

```c
typedef struct {
    int32 length;
    char data[FLEXIBLE_ARRAY_MEMBER];
} text;
```

The \texttt{[FLEXIBLE_ARRAY_MEMBER]} notation means that the actual length of the data part is not specified by this declaration.

When manipulating variable-length types, we must be careful to allocate the correct amount of memory and set the length field correctly. For example, if we wanted to store 40 bytes in a \texttt{text} structure, we might use a code fragment like this:

```c
#include "postgres.h"
...
char buffer[40]; /* our source data */
...
text *destination = (text *) palloc(VARHDRSZ + 40);
SET_VARSIZE(destination, VARHDRSZ + 40);
```
memcpy(destination->data, buffer, 40);
...

VARHDRSZ is the same as sizeof(int32), but it’s considered good style to use the macro VARHDRSZ to refer to the size of the overhead for a variable-length type. Also, the length field must be set using the SET_VARSIZE macro, not by simple assignment.

Table 36-1 specifies which C type corresponds to which SQL type when writing a C-language function that uses a built-in type of PostgreSQL. The “Defined In” column gives the header file that needs to be included to get the type definition. (The actual definition might be in a different file that is included by the listed file. It is recommended that users stick to the defined interface.) Note that you should always include postgres.h first in any source file, because it declares a number of things that you will need anyway.

<table>
<thead>
<tr>
<th>SQL Type</th>
<th>C Type</th>
<th>Defined In</th>
</tr>
</thead>
<tbody>
<tr>
<td>abstime</td>
<td>AbsoluteTime</td>
<td>utils/nabstime.h</td>
</tr>
<tr>
<td>bigint(int8)</td>
<td>int64</td>
<td>postgres.h</td>
</tr>
<tr>
<td>boolean</td>
<td>bool</td>
<td>postgres.h (maybe compiler</td>
</tr>
<tr>
<td>box</td>
<td>BOX*</td>
<td>utils/geo_decls.h</td>
</tr>
<tr>
<td>bytea</td>
<td>bytea*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>*char</td>
<td>char</td>
<td>(compiler built-in)</td>
</tr>
<tr>
<td>character</td>
<td>BpChar*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>cid</td>
<td>CommandId</td>
<td>postgres.h</td>
</tr>
<tr>
<td>date</td>
<td>DateADT</td>
<td>utils/date.h</td>
</tr>
<tr>
<td>smallint(int2)</td>
<td>int16</td>
<td>postgres.h</td>
</tr>
<tr>
<td>int2vector</td>
<td>int2vector*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>integer(int4)</td>
<td>int32</td>
<td>postgres.h</td>
</tr>
<tr>
<td>real(float4)</td>
<td>float4*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>double precision(float8)</td>
<td>float8*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>interval</td>
<td>Interval*</td>
<td>datatype/timestamp.h</td>
</tr>
<tr>
<td>lseg</td>
<td>LSEG*</td>
<td>utils/geo_decls.h</td>
</tr>
<tr>
<td>name</td>
<td>Name</td>
<td>postgres.h</td>
</tr>
<tr>
<td>oid</td>
<td>Oid</td>
<td>postgres.h</td>
</tr>
<tr>
<td>oidvector</td>
<td>oidvector*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>path</td>
<td>PATH*</td>
<td>utils/geo_decls.h</td>
</tr>
<tr>
<td>point</td>
<td>POINT*</td>
<td>utils/geo_decls.h</td>
</tr>
<tr>
<td>regproc</td>
<td>regproc</td>
<td>postgres.h</td>
</tr>
<tr>
<td>reltime</td>
<td>RelativeTime</td>
<td>utils/nabstime.h</td>
</tr>
<tr>
<td>text</td>
<td>text*</td>
<td>postgres.h</td>
</tr>
<tr>
<td>tid</td>
<td>ItemPointer</td>
<td>storage/itemptptr.h</td>
</tr>
<tr>
<td>time</td>
<td>TimeADT</td>
<td>utils/date.h</td>
</tr>
<tr>
<td>time with time zone</td>
<td>TimeTzADT</td>
<td>utils/date.h</td>
</tr>
<tr>
<td>timestamp</td>
<td>Timestamp*</td>
<td>datatype/timestamp.h</td>
</tr>
</tbody>
</table>
### 36.9.3. Version 0 Calling Conventions

We present the “old style” calling convention first — although this approach is now deprecated, it’s easier to get a handle on initially. In the version-0 method, the arguments and result of the C function are just declared in normal C style, but being careful to use the C representation of each SQL data type as shown above.

Here are some examples:

```c
#include "postgres.h"
#include <string.h>
#include "utils/geo_decls.h"

#ifdef PG_MODULE_MAGIC
PG_MODULE_MAGIC;
#endif

/* by value */

int add_one(int arg)
{
    return arg + 1;
}

/* by reference, fixed length */

float8 *
add_one_float8(float8 *arg)
{
    float8  *result = (float8 *) palloc(sizeof(float8));
    *result = *arg + 1.0;
    return result;
}

Point *
makepoint(Point *pointx, Point *pointy)
{
    Point  *new_point = (Point *) palloc(sizeof(Point));
    new_point->x = pointx->x;
    new_point->y = pointy->y;
    return new_point;
}
```
/* by reference, variable length */

text *
copytext(text *t)
{
    /*
    * VARSIZE is the total size of the struct in bytes.
    */
text *new_t = (text *) palloc(VARSIZE(t));
SET_VARSIZE(new_t, VARSIZE(t));
    /*
    * VARDATA is a pointer to the data region of the struct.
    */
copy((void *) VARDATA(new_t), /* destination */
    (void *) VARDATA(t), /* source */
    VARSIZE(t) - VARHDRSZ); /* how many bytes */
return new_t;
}

text *
concat_text(text *arg1, text *arg2)
{
    int32 new_text_size = VARSIZE(arg1) + VARSIZE(arg2) - VARHDRSZ;
text *new_text = (text *) palloc(new_text_size);

    SET_VARSIZE(new_text, new_text_size);
copy((void *) VARDATA(new_text), VARDATA(arg1), VARSIZE(arg1) - VARHDRSZ);
copy((void *) VARDATA(new_text) + (VARSIZE(arg1) - VARHDRSZ),
    VARDATA(arg2), VARSIZE(arg2) - VARHDRSZ);
return new_text;
}

Supposing that the above code has been prepared in file funcs.c and compiled into a shared object, we could define the functions to PostgreSQL with commands like this:

CREATE FUNCTION add_one(integer) RETURNS integer
    AS 'DIRECTORY/funcs', 'add_one'
    LANGUAGE C STRICT;

-- note overloading of SQL function name "add_one"
CREATE FUNCTION add_one(double precision) RETURNS double precision
    AS 'DIRECTORY/funcs', 'add_one_float8'
    LANGUAGE C STRICT;

CREATE FUNCTION makepoint(point, point) RETURNS point
    AS 'DIRECTORY/funcs', 'makepoint'
    LANGUAGE C STRICT;

CREATE FUNCTION copytext(text) RETURNS text
    AS 'DIRECTORY/funcs', 'copytext'
    LANGUAGE C STRICT;

CREATE FUNCTION concat_text(text, text) RETURNS text
    AS 'DIRECTORY/funcs', 'concat_text'
    LANGUAGE C STRICT;
Here, DIRECTORY stands for the directory of the shared library file (for instance the PostgreSQL tutorial directory, which contains the code for the examples used in this section). (Better style would be to use just 'funcn' in the AS clause, after having added DIRECTORY to the search path. In any case, we can omit the system-specific extension for a shared library, commonly .so or .sl.)

Notice that we have specified the functions as “strict”, meaning that the system should automatically assume a null result if any input value is null. By doing this, we avoid having to check for null inputs in the function code. Without this, we’d have to check for null values explicitly, by checking for a null pointer for each pass-by-reference argument. (For pass-by-value arguments, we don’t even have a way to check!)

Although this calling convention is simple to use, it is not very portable; on some architectures there are problems with passing data types that are smaller than int this way. Also, there is no simple way to return a null result, nor to cope with null arguments in any way other than making the function strict. The version-1 convention, presented next, overcomes these objections.

### 36.9.4. Version 1 Calling Conventions

The version-1 calling convention relies on macros to suppress most of the complexity of passing arguments and results. The C declaration of a version-1 function is always:

```c
Datum funcname(PG_FUNCTION_ARGS)
```

In addition, the macro call:

```c
PG_FUNCTION_INFO_V1(funcname);
```

must appear in the same source file. (Conventionally, it’s written just before the function itself.) This macro call is not needed for internal-language functions, since PostgreSQL assumes that all internal functions use the version-1 convention. It is, however, required for dynamically-loaded functions. In a version-1 function, each actual argument is fetched using a PG_GETARG_xxx() macro that corresponds to the argument’s data type, and the result is returned using a PG_RETURN_xxx() macro for the return type. PG_GETARG_xxx() takes as its argument the number of the function argument to fetch, where the count starts at 0. PG_RETURN_xxx() takes as its argument the actual value to return.

Here we show the same functions as above, coded in version-1 style:

```c
#include "postgres.h"
#include <string.h>
#include "fmgr.h"
#include "utils/geo_decls.h"

#ifdef PG_MODULE_MAGIC
PG_MODULE_MAGIC;
#endif

/* by value */

PG_FUNCTION_INFO_V1(add_one);

Datum add_one(PG_FUNCTION_ARGS)
{
    int32 arg = PG_GETARG_INT32(0);
    return arg + 1;
}
```
Chapter 36. Extending SQL

```c
PG_RETURN_INT32(arg + 1);
}
/* by reference, fixed length */
PG_FUNCTION_INFO_V1(add_one_float8);
Datum
add_one_float8(PG_FUNCTION_ARGS)
{
    /* The macros for FLOAT8 hide its pass-by-reference nature. */
    float8 arg = PG_GETARG_FLOAT8(0);
    PG_RETURN_FLOAT8(arg + 1.0);
}

PG_FUNCTION_INFO_V1(makepoint);
Datum
makepoint(PG_FUNCTION_ARGS)
{
    /* Here, the pass-by-reference nature of Point is not hidden. */
    Point *pointx = PG_GETARG_POINT_P(0);
    Point *pointy = PG_GETARG_POINT_P(1);
    Point *new_point = (Point *) palloc(sizeof(Point));
    new_point->x = pointx->x;
    new_point->y = pointy->y;
    PG_RETURN_POINT_P(new_point);
}
/* by reference, variable length */
PG_FUNCTION_INFO_V1(copytext);
Datum
copytext(PG_FUNCTION_ARGS)
{
    text *t = PG_GETARG_TEXT_P(0);
    /*
        * VARSIZE is the total size of the struct in bytes.
        */
    text *new_t = (text *) palloc(VARSIZE(t));
    SET_VARSIZE(new_t, VARSIZE(t));
    /*
        * VARDATA is a pointer to the data region of the struct.
        */
    memcpy((void *) VARDATA(new_t), /* destination */
            (void *) VARDATA(t), /* source */
            VARSIZE(t) - VARHDRSZ); /* how many bytes */
    PG_RETURN_TEXT_P(new_t);
}
PG_FUNCTION_INFO_V1(concat_text);
```
Datum
concat_text(PG_FUNCTION_ARGS)
{
  text *arg1 = PG_GETARG_TEXT_P(0);
text *arg2 = PG_GETARG_TEXT_P(1);
int32 new_text_size = VARSIZE(arg1) + VARSIZE(arg2) - VARHDRSZ;
text *new_text = (text *) palloc(new_text_size);
 SET_VARSIZE(new_text, new_text_size);
 memcpy(VARDATA(new_text), VARDATA(arg1), VARSIZE(arg1) - VARHDRSZ);
 memcpy(VARDATA(new_text) + (VARSIZE(arg1) - VARHDRSZ),
         VARDATA(arg2), VARSIZE(arg2) - VARHDRSZ);
 PG_RETURN_TEXT_P(new_text);
}

The CREATE FUNCTION commands are the same as for the version-0 equivalents.

At first glance, the version-1 coding conventions might appear to be just pointless obscurantism. They
do, however, offer a number of improvements, because the macros can hide unnecessary detail. An
example is that in coding add_one_float8, we no longer need to be aware that float8 is a pass-
by-reference type. Another example is that the GETARG macros for variable-length types allow for
more efficient fetching of “toasted” (compressed or out-of-line) values.

One big improvement in version-1 functions is better handling of null inputs and results. The macro
PG_ARGISNULL(n) allows a function to test whether each input is null. (Of course, doing this is
only necessary in functions not declared “strict.”) As with the PG_GETARG_xxx() macros, the
input arguments are counted beginning at zero. Note that one should refrain from executing
PG_GETARG_xxx() until one has verified that the argument isn’t null. To return a null result, execute
PG_RETURN_NULL(); this works in both strict and nonstrict functions.

Other options provided in the new-style interface are two variants of the PG_GETARG_xxx() macros.
The first of these, PG_GETARG_xxx_COPY(), guarantees to return a copy of the specified argument
that is safe for writing into. (The normal macros will sometimes return a pointer to a value that is
physically stored in a table, which must not be written to. Using the PG_GETARG_xxx_COPY() macros
guarantees a writable result.) The second variant consists of the PG_GETARG_xxx_SLICE() macros
which take three arguments. The first is the number of the function argument (as above). The
second and third are the offset and length of the segment to be returned. Offsets are counted from
zero, and a negative length requests that the remainder of the value be returned. These macros provide
more efficient access to parts of large values in the case where they have storage type “external”.
(The storage type of a column can be specified using ALTER TABLE tablename ALTER COLUMN
colname SET STORAGE storagetype. storagetype is one of plain, external, extended, or
main.)

Finally, the version-1 function call conventions make it possible to return set results (Section 36.9.9)
and implement trigger functions (Chapter 37) and procedural-language call handlers (Chapter 54).
Version-1 code is also more portable than version-0, because it does not break restrictions on function
call protocol in the C standard. For more details see src/backend/utils/fmgr/README in the
source distribution.

36.9.5. Writing Code

Before we turn to the more advanced topics, we should discuss some coding rules for PostgreSQL
C-language functions. While it might be possible to load functions written in languages other than
C into PostgreSQL, this is usually difficult (when it is possible at all) because other languages, such as C++, FORTRAN, or Pascal often do not follow the same calling convention as C. That is, other languages do not pass argument and return values between functions in the same way. For this reason, we will assume that your C-language functions are actually written in C.

The basic rules for writing and building C functions are as follows:

- Use `pg_config --includedir-server` to find out where the PostgreSQL server header files are installed on your system (or the system that your users will be running on).
- Compiling and linking your code so that it can be dynamically loaded into PostgreSQL always requires special flags. See Section 36.9.6 for a detailed explanation of how to do it for your particular operating system.
- Remember to define a “magic block” for your shared library, as described in Section 36.9.1.
- When allocating memory, use the PostgreSQL functions `palloc` and `pfree` instead of the corresponding C library functions `malloc` and `free`. The memory allocated by `palloc` will be freed automatically at the end of each transaction, preventing memory leaks.
- Always zero the bytes of your structures using `memset` (or allocate them with `palloc0` in the first place). Even if you assign to each field of your structure, there might be alignment padding (holes in the structure) that contain garbage values. Without this, it’s difficult to support hash indexes or hash joins, as you must pick out only the significant bits of your data structure to compute a hash. The planner also sometimes relies on comparing constants via bitwise equality, so you can get undesirable planning results if logically-equivalent values aren’t bitwise equal.
- Most of the internal PostgreSQL types are declared in `postgres.h`, while the function manager interfaces (PG_FUNCTION_ARGS, etc.) are in `fmgr.h`, so you will need to include at least these two files. For portability reasons it’s best to include `postgres.h` first, before any other system or user header files. Including `postgres.h` will also include `elog.h` and `palloc.h` for you.
- Symbol names defined within object files must not conflict with each other or with symbols defined in the PostgreSQL server executable. You will have to rename your functions or variables if you get error messages to this effect.

### 36.9.6. Compiling and Linking Dynamically-loaded Functions

Before you are able to use your PostgreSQL extension functions written in C, they must be compiled and linked in a special way to produce a file that can be dynamically loaded by the server. To be precise, a shared library needs to be created.

For information beyond what is contained in this section you should read the documentation of your operating system, in particular the manual pages for the C compiler, `cc`, and the link editor, `ld`. In addition, the PostgreSQL source code contains several working examples in the `contrib` directory. If you rely on these examples you will make your modules dependent on the availability of the PostgreSQL source code, however.

Creating shared libraries is generally analogous to linking executables: first the source files are compiled into object files, then the object files are linked together. The object files need to be created as position-independent code (PIC), which conceptually means that they can be placed at an arbitrary location in memory when they are loaded by the executable. (Object files intended for executables are usually not compiled that way.) The command to link a shared library contains special flags to
distinguish it from linking an executable (at least in theory — on some systems the practice is much uglier).

In the following examples we assume that your source code is in a file foo.c and we will create a shared library foo.so. The intermediate object file will be called foo.o unless otherwise noted. A shared library can contain more than one object file, but we only use one here.

FreeBSD

The compiler flag to create PIC is -fPIC. To create shared libraries the compiler flag is -shared.

```
 gcc -fPIC -c foo.c
 gcc -shared -o foo.so foo.o
```

This is applicable as of version 3.0 of FreeBSD.

HP-UX

The compiler flag of the system compiler to create PIC is +z. When using GCC it's -fPIC. The linker flag for shared libraries is -b. So:

```
 cc +z -c foo.c
 or:
 gcc -fPIC -c foo.c
 and then:
 ld -b -o foo.sl foo.o
```

HP-UX uses the extension .sl for shared libraries, unlike most other systems.

Linux

The compiler flag to create PIC is -fPIC. The compiler flag to create a shared library is -shared. A complete example looks like this:

```
 cc -fPIC -c foo.c
 cc -shared -o foo.so foo.o
```

OS X

Here is an example. It assumes the developer tools are installed.

```
 cc -c foo.c
 cc -bundle -flat_namespace -undefined suppress -o foo.so foo.o
```

NetBSD

The compiler flag to create PIC is -fPIC. For ELF systems, the compiler with the flag -shared is used to link shared libraries. On the older non-ELF systems, ld -Bshareable is used.

```
 gcc -fPIC -c foo.c
 gcc -shared -o foo.so foo.o
```

OpenBSD

The compiler flag to create PIC is -fPIC. ld -Bshareable is used to link shared libraries.

```
 gcc -fPIC -c foo.c
 ld -Bshareable -o foo.so foo.o
```

Solaris

The compiler flag to create PIC is -KPIC with the Sun compiler and -fPIC with GCC. To link shared libraries, the compiler option is -G with either compiler or alternatively -shared with GCC.

```
 cc -KPIC -c foo.c
 cc -G -o foo.so foo.o
```
Chapter 36. Extending SQL

UnixWare

The compiler flag to create PIC is -K PIC with the SCO compiler and -fpic with GCC. To link shared libraries, the compiler option is -G with the SCO compiler and -shared with GCC.

```
cc -K PIC -c foo.c
cc -G -o foo.so foo.o
```

```
gcc -fpic -c foo.c
gcc -shared -o foo.so foo.o
```

**Tip:** If this is too complicated for you, you should consider using GNU Libtool¹, which hides the platform differences behind a uniform interface.

The resulting shared library file can then be loaded into PostgreSQL. When specifying the file name to the `CREATE FUNCTION` command, one must give it the name of the shared library file, not the intermediate object file. Note that the system’s standard shared-library extension (usually `.so` or `.sl`) can be omitted from the `CREATE FUNCTION` command, and normally should be omitted for best portability.

Refer back to Section 36.9.1 about where the server expects to find the shared library files.

### 36.9.7. Composite-type Arguments

Composite types do not have a fixed layout like C structures. Instances of a composite type can contain null fields. In addition, composite types that are part of an inheritance hierarchy can have different fields than other members of the same inheritance hierarchy. Therefore, PostgreSQL provides a function interface for accessing fields of composite types from C.

Suppose we want to write a function to answer the query:

```
SELECT name, c_overpaid(emp, 1500) AS overpaid
    FROM emp
    WHERE name = 'Bill' OR name = 'Sam';
```

Using call conventions version 0, we can define `c_overpaid` as:

```
#include "postgres.h"
#include "executor/executor.h" /* for GetAttributeByName() */

#ifdef PG_MODULE_MAGIC
PG_MODULE_MAGIC;
#endif

bool
c_overpaid(HeapTupleHeader t, /* the current row of emp */
            int32 limit)
{
    bool isnull;
```

int32 salary;

salary = DatumGetInt32(GetAttributeByName(t, "salary", &isnull));
if (isnull)
    return false;
return salary > limit;
}

In version-1 coding, the above would look like this:

#include "postgres.h"
#include "executor/executor.h" /* for GetAttributeByName() */

#ifdef PG_MODULE_MAGIC
PG_MODULE_MAGIC;
#endif

PG_FUNCTION_INFO_V1(c_overpaid);

Datum

c_overpaid(PG_FUNCTION_ARGS)
{
    HeapTupleHeader t = PG_GETARG_HEAPTUPLEHEADER(0);
    int32 limit = PG_GETARG_INT32(1);
    bool isnull;
    Datum salary;

    salary = GetAttributeByName(t, "salary", &isnull);
    if (isnull)
        PG_RETURN_BOOL(false);
    /* Alternatively, we might prefer to do PG_RETURN_NULL() for null salary. */
    PG_RETURN_BOOL(DatumGetInt32(salary) > limit);
}

GetAttributeByName is the PostgreSQL system function that returns attributes out of the specified row. It has three arguments: the argument of type HeapTupleHeader passed into the function, the name of the desired attribute, and a return parameter that tells whether the attribute is null. GetAttributeByName returns a Datum value that you can convert to the proper data type by using the appropriate DatumGetXXX() macro. Note that the return value is meaningless if the null flag is set; always check the null flag before trying to do anything with the result.

There is also GetAttributeByNum, which selects the target attribute by column number instead of name.

The following command declares the function c_overpaid in SQL:

CREATE FUNCTION c_overpaid(emp, integer) RETURNS boolean
    AS 'DIRECTORY/funcs', 'c_overpaid'
    LANGUAGE C STRICT;

Notice we have used STRICT so that we did not have to check whether the input arguments were NULL.
36.9.8. Returning Rows (Composite Types)

To return a row or composite-type value from a C-language function, you can use a special API that provides macros and functions to hide most of the complexity of building composite data types. To use this API, the source file must include:

```c
#include "funcapi.h"
```

There are two ways you can build a composite data value (henceforth a “tuple”): you can build it from an array of Datum values, or from an array of C strings that can be passed to the input conversion functions of the tuple’s column data types. In either case, you first need to obtain or construct a `TupleDesc` descriptor for the tuple structure. When working with Datums, you pass the `TupleDesc` to `BlessTupleDesc`, and then call `heap_form_tuple` for each row. When working with C strings, you pass the `TupleDesc` to `TupleDescGetAttInMetadata`, and then call `BuildTupleFromCStrings` for each row. In the case of a function returning a set of tuples, the setup steps can all be done once during the first call of the function.

Several helper functions are available for setting up the needed `TupleDesc`. The recommended way to do this in most functions returning composite values is to call:

```c
TypeFuncClass get_call_result_type(FunctionCallInfo fcinfo,
                                   Oid *resultTypeId,
                                   TupleDesc *resultTupleDesc)
```

passing the same `fcinfo` struct passed to the calling function itself. (This of course requires that you use the version-1 calling conventions.) `resultTypeId` can be specified as `NULL` or as the address of a local variable to receive the function’s result type OID. `resultTupleDesc` should be the address of a local `TupleDesc` variable. Check that the result is `TYPEFUNC_COMPOSITE`; if so, `resultTupleDesc` has been filled with the needed `TupleDesc`. (If it is not, you can report an error along the lines of “function returning record called in context that cannot accept type record”.)

**Tip:** `get_call_result_type` can resolve the actual type of a polymorphic function result; so it is useful in functions that return scalar polymorphic results, not only functions that return composites. The `resultTypeId` output is primarily useful for functions returning polymorphic scalars.

**Note:** `get_call_result_type` has a sibling `get_expr_result_type`, which can be used to resolve the expected output type for a function call represented by an expression tree. This can be used when trying to determine the result type from outside the function itself. There is also `get_func_result_type`, which can be used when only the function’s OID is available. However, these functions are not able to deal with functions declared to return record, and `get_func_result_type` cannot resolve polymorphic types, so you should preferentially use `get_call_result_type`.

Older, now-deprecated functions for obtaining `TupleDescs` are:

```c
TupleDesc RelationNameGetTupleDesc(const char *relname)
```

to get a `TupleDesc` for the row type of a named relation, and:

```c
TupleDesc TypeGetTupleDesc(Oid typeoid, List *colaliases)
```

1017
to get a TupleDesc based on a type OID. This can be used to get a TupleDesc for a base or composite type. It will not work for a function that returns record, however, and it cannot resolve polymorphic types.

Once you have a TupleDesc, call:

TupleDesc BlessTupleDesc(TupleDesc tupdesc)

if you plan to work with Datums, or:

AttInMetadata *TupleDescGetAttInMetadata(TupleDesc tupdesc)

if you plan to work with C strings. If you are writing a function returning set, you can save the results of these functions in the FuncCallContext structure — use the tuple_desc or attinmeta field respectively.

When working with Datums, use:

HeapTuple heap_form_tuple(TupleDesc tupdesc, Datum *values, bool *isnull)

to build a HeapTuple given user data in Datum form.

When working with C strings, use:

HeapTuple BuildTupleFromCStrings(AttInMetadata *attinmeta, char **values)

to build a HeapTuple given user data in C string form. values is an array of C strings, one for each attribute of the return row. Each C string should be in the form expected by the input function of the attribute data type. In order to return a null value for one of the attributes, the corresponding pointer in the values array should be set to NULL. This function will need to be called again for each row you return.

Once you have built a tuple to return from your function, it must be converted into a Datum. Use:

HeapTupleGetDatum(HeapTuple tuple)

to convert a HeapTuple into a valid Datum. This Datum can be returned directly if you intend to return just a single row, or it can be used as the current return value in a set-returning function.

An example appears in the next section.

36.9.9. Returning Sets

There is also a special API that provides support for returning sets (multiple rows) from a C-language function. A set-returning function must follow the version-1 calling conventions. Also, source files must include funcapi.h, as above.

A set-returning function (SRF) is called once for each item it returns. The SRF must therefore save enough state to remember what it was doing and return the next item on each call. The structure FuncCallContext is provided to help control this process. Within a function, fcinfo->flinfo->fn_extra is used to hold a pointer to FuncCallContext across calls.

typedef struct FuncCallContext
{
    /*
        * Number of times we’ve been called before
        *
        * call_cntr is initialized to 0 for you by SRF_FIRSTCALL_INIT(), and
    */
Chapter 36. Extending SQL

* incremented for you every time SRF_RETURN_NEXT() is called.
* 
uint64 call_cntr;

/*
* OPTIONAL maximum number of calls
* max_calls is here for convenience only and setting it is optional.
* If not set, you must provide alternative means to know when the
* function is done.
*/
uint64 max_calls;

/*
* OPTIONAL pointer to result slot
* This is obsolete and only present for backward compatibility, viz,
* user-defined SRFs that use the deprecated TupleDescGetSlot().
*/
TupleTableSlot *slot;

/*
* OPTIONAL pointer to miscellaneous user-provided context information
* user_fctx is for use as a pointer to your own data to retain
* arbitrary context information between calls of your function.
*/
void *user_fctx;

/*
* OPTIONAL pointer to struct containing attribute type input metadata
* attinmeta is for use when returning tuples (i.e., composite data types)
* and is not used when returning base data types. It is only needed
* if you intend to use BuildTupleFromCStrings() to create the return
* tuple.
*/
AttInMetadata *attinmeta;

/*
* memory context used for structures that must live for multiple calls
* multi_call_memory_ctx is set by SRF_FIRSTCALL_INIT() for you, and used
* by SRF_RETURN_DONE() for cleanup. It is the most appropriate memory
* context for any memory that is to be reused across multiple calls
* of the SRF.
*/
MemoryContext multi_call_memory_ctx;

/*
* OPTIONAL pointer to struct containing tuple description
* tuple_desc is for use when returning tuples (i.e., composite data types)
* and is only needed if you are going to build the tuples with
* heap_form_tuple() rather than with BuildTupleFromCStrings(). Note that
* the TupleDesc pointer stored here should usually have been run through
* BlessTupleDesc() first.
*/
An SRF uses several functions and macros that automatically manipulate the `FuncCallContext` structure (and expect to find it via `fn_extra`). Use:

```
SRF_IS_FIRSTCALL()
```

to determine if your function is being called for the first or a subsequent time. On the first call (only) use:

```
SRF_FIRSTCALL_INIT()
```

to initialize the `FuncCallContext`. On every function call, including the first, use:

```
SRF_PERCALL_SETUP()
```

to properly set up for using the `FuncCallContext` and clearing any previously returned data left over from the previous pass.

If your function has data to return, use:

```
SRF_RETURN_NEXT(funcctx, result)
```

to return it to the caller. (result must be of type Datum, either a single value or a tuple prepared as described above.) Finally, when your function is finished returning data, use:

```
SRF_RETURN_DONE(funcctx)
```

to clean up and end the SRF.

The memory context that is current when the SRF is called is a transient context that will be cleared between calls. This means that you do not need to call `pfree` on everything you allocated using `palloc`; it will go away anyway. However, if you want to allocate any data structures to live across calls, you need to put them somewhere else. The memory context referenced by `multi_call_memory_ctx` is a suitable location for any data that needs to survive until the SRF is finished running. In most cases, this means that you should switch into `multi_call_memory_ctx` while doing the first-call setup.

**Warning**

While the actual arguments to the function remain unchanged between calls, if you detoast the argument values (which is normally done transparently by the `PG_GETARG_` macro) in the transient context then the detoasted copies will be freed on each cycle. Accordingly, if you keep references to such values in your `user_fctx`, you must either copy them into the `multi_call_memory_ctx` after detoasting, or ensure that you detoast the values only in that context.

A complete pseudo-code example looks like the following:

```c
Datum
my_set_returning_function(PG_FUNCTION_ARGS)
{
    FuncCallContext  *funcctx;
    Datum    result;
    /*
       TupleDesc tuple_desc;
    */
```
Chapter 36. Extending SQL

further declarations as needed

if (SRF_IS_FIRSTCALL())
{
    MemoryContext oldcontext;

    funcctx = SRF_FIRSTCALL_INIT();
    oldcontext = MemoryContextSwitchTo(funcctx->multi_call_memory_ctx);
    /* One-time setup code appears here: */
    user code
    if returning composite
        build TupleDesc, and perhaps AttInMetadata
    endif returning composite
    user code
    MemoryContextSwitchTo(oldcontext);
}

/* Each-time setup code appears here: */
user code
funcctx = SRF_PERCALL_SETUP();
user code

/* this is just one way we might test whether we are done: */
if (funcctx->call_cntr < funcctx->max_calls)
{
    /* Here we want to return another item: */
    user code
    obtain result Datum
    SRF_RETURN_NEXT(funcctx, result);
}
else
{
    /* Here we are done returning items and just need to clean up: */
    user code
    SRF_RETURN_DONE(funcctx);
}
}

A complete example of a simple SRF returning a composite type looks like:

PG_FUNCTION_INFO_V1(retcomposite);

Datum retcomposite(PG_FUNCTION_ARGS)
{
    FuncCallContext *funcctx;
    int call_cntr;
    int max_calls;
    TupleDesc tupdesc;
    AttInMetadata *attinmeta;

    /* stuff done only on the first call of the function */
    if (SRF_IS_FIRSTCALL())
    {
        MemoryContext oldcontext;

        funcctx = SRF_FIRSTCALL_INIT();
        oldcontext = MemoryContextSwitchTo(funcctx->multi_call_memory_ctx);
    
        /* One-time setup code appears here: */
        user code
        if returning composite
            build TupleDesc, and perhaps AttInMetadata
        endif returning composite
        user code
        MemoryContextSwitchTo(oldcontext);
    }
Chapter 36. Extending SQL

/* create a function context for cross-call persistence */
funcctx = SRF_FIRSTCALL_INIT();

/* switch to memory context appropriate for multiple function calls */
oldcontext = MemoryContextSwitchTo(funcctx->multi_call_memory_ctx);

/* total number of tuples to be returned */
funcctx->max_calls = PG_GETARG_UINT32(0);

/* Build a tuple descriptor for our result type */
if (get_call_result_type(fcinfo, NULL, &tupdesc) != TYPEFUNC_COMPOSITE)
ereport(ERROR,
{errcode(ERRCODE_FEATURE_NOT_SUPPORTED),
 errmsg("function returning record called in context "
"that cannot accept type record")});

/* generate attribute metadata needed later to produce tuples from raw
* C strings */
attinmeta = TupleDescGetAttInMetadata(tupdesc);
funcctx->attinmeta = attinmeta;
MemoryContextSwitchTo(oldcontext);

} /* stuff done on every call of the function */
funcctx = SRF_PERCALL_SETUP();
call_cntr = funcctx->call_cntr;
max_calls = funcctx->max_calls;
attinmeta = funcctx->attinmeta;

if (call_cntr < max_calls) /* do when there is more left to send */
{
char **values;
HeapTuple tuple;
Datum result;

/* Prepare a values array for building the returned tuple.
* This should be an array of C strings which will
* be processed later by the type input functions.
*/
values = (char **) palloc(3 * sizeof(char *));
values[0] = (char *) palloc(16 * sizeof(char));
values[1] = (char *) palloc(16 * sizeof(char));
values[2] = (char *) palloc(16 * sizeof(char));

snprintf(values[0], 16, "%d", 1 * PG_GETARG_INT32(1));
snprintf(values[1], 16, "%d", 2 * PG_GETARG_INT32(1));
snprintf(values[2], 16, "%d", 3 * PG_GETARG_INT32(1));

/* build a tuple */
tuple = BuildTupleFromCStrings(attinmeta, values);

/* make the tuple into a datum */

1022
One way to declare this function in SQL is:

```sql
CREATE TYPE __retcomposite AS (f1 integer, f2 integer, f3 integer);
CREATE OR REPLACE FUNCTION retcomposite(integer, integer)
RETURNS SETOF __retcomposite
AS 'filename', 'retcomposite'
LANGUAGE C IMMUTABLE STRICT;
```

A different way is to use OUT parameters:

```sql
CREATE OR REPLACE FUNCTION retcomposite(IN integer, IN integer,
OUT f1 integer, OUT f2 integer, OUT f3 integer)
RETURNS SETOF record
AS 'filename', 'retcomposite'
LANGUAGE C IMMUTABLE STRICT;
```

Notice that in this method the output type of the function is formally an anonymous record type.

The directory contrib/tablefunc module in the source distribution contains more examples of set-returning functions.

### 36.9.10. Polymorphic Arguments and Return Types

C-language functions can be declared to accept and return the polymorphic types anyelement, anyarray, anynonarray, anyenum, and anyrange. See Section 36.2.5 for a more detailed explanation of polymorphic functions. When function arguments or return types are defined as polymorphic types, the function author cannot know in advance what data type it will be called with, or need to return. There are two routines provided in fmgr.h to allow a version-1 C function to discover the actual data types of its arguments and the type it is expected to return. The routines are called `get_fn_expr_rettype(FmgrInfo *flinfo)` and `get_fn_expr_argtype(FmgrInfo *flinfo, int argnum)`. They return the result or argument type OID, or InvalidOid if the information is not available. The structure flinfo is normally accessed as fcinfo->flinfo. The parameter argnum is zero based. `get_call_result_type` can also be used as an alternative to `get_fn_expr_rettype`. There is also `get_fn_expr_variadic`, which can be used to find out whether variadic arguments have been merged into an array. This is primarily useful for VARIADIC "any" functions, since such merging will always have occurred for variadic functions taking ordinary array types.
For example, suppose we want to write a function to accept a single element of any type, and return a one-dimensional array of that type:

```c
PG_FUNCTION_INFO_V1(make_array);
Datum make_array(PG_FUNCTION_ARGS)
{
    ArrayType *result;
    Oid element_type = get_fn_expr_argtype(fcinfo->flinfo, 0);
    Datum element;
    bool isnull;
    int16 typlen;
    bool typbyval;
    char typalign;
    int ndims;
    int dims[MAXDIM];
    int lbs[MAXDIM];

    if (!OidIsValid(element_type))
        elog(ERROR, "could not determine data type of input");

    /* get the provided element, being careful in case it’s NULL */
    isnull = PG_ARGISNULL(0);
    if (isnull)
        element = (Datum) 0;
    else
        element = PG_GETARG_DATUM(0);

    /* we have one dimension */
    ndims = 1;
    /* and one element */
    dims[0] = 1;
    /* and lower bound is 1 */
    lbs[0] = 1;

    /* get required info about the element type */
    get_typlenbyvalalign(element_type, &typlen, &typbyval, &typalign);

    /* now build the array */
    result = construct_md_array(&element, &isnull, ndims, dims, lbs,
                                element_type, typlen, typbyval, typalign);

    PG_RETURN_ARRAYTYPE_P(result);
}
```

The following command declares the function `make_array` in SQL:

```sql
CREATE FUNCTION make_array(anyelement) RETURNS anyarray
AS 'DIRECTORY/funcs', 'make_array'
LANGUAGE C IMMUTABLE;
```

There is a variant of polymorphism that is only available to C-language functions: they can be declared to take parameters of type "any". (Note that this type name must be double-quoted, since it’s also a SQL reserved word.) This works like `anyelement` except that it does not constrain different "any"
arguments to be the same type, nor do they help determine the function’s result type. A C-language function can also declare its final parameter to be `VARIADIC "any"`. This will match one or more actual arguments of any type (not necessarily the same type). These arguments will not be gathered into an array as happens with normal variadic functions; they will just be passed to the function separately. The `PG_NARGS()` macro and the methods described above must be used to determine the number of actual arguments and their types when using this feature. Also, users of such a function might wish to use the `VARIADIC` keyword in their function call, with the expectation that the function would treat the array elements as separate arguments. The function itself must implement that behavior if wanted, after using `get_fn_expr_variadic` to detect that the actual argument was marked with `VARIADIC`.

### 36.9.11. Transform Functions

Some function calls can be simplified during planning based on properties specific to the function. For example, `int4mul(n, 1)` could be simplified to just `n`. To define such function-specific optimizations, write a `transform function` and place its OID in the `protransform` field of the primary function’s `pg_proc` entry. The transform function must have the SQL signature `protransform(internal) RETURNS internal`. The argument, actually `FuncExpr *`, is a dummy node representing a call to the primary function. If the transform function’s study of the expression tree proves that a simplified expression tree can substitute for all possible concrete calls represented thereby, build and return that simplified expression. Otherwise, return a `NULL` pointer (not a SQL null).

We make no guarantee that PostgreSQL will never call the primary function in cases that the transform function could simplify. Ensure rigorous equivalence between the simplified expression and an actual call to the primary function.

Currently, this facility is not exposed to users at the SQL level because of security concerns, so it is only practical to use for optimizing built-in functions.

### 36.9.12. Shared Memory and LWLocks

Add-ins can reserve LWLocks and an allocation of shared memory on server startup. The add-in’s shared library must be preloaded by specifying it in `shared_preload_libraries`. Shared memory is reserved by calling:

```c
void RequestAddinShmemSpace(int size)
```

from your `_PG_init` function.

LWLocks are reserved by calling:

```c
void RequestNamedLWLockTranche(const char *tranche_name, int num_lwlocks)
```

from `_PG_init`. This will ensure that an array of `num_lwlocks` LWLocks is available under the name `tranche_name`. Use `GetNamedLWLockTranche` to get a pointer to this array.

To avoid possible race-conditions, each backend should use the LWLock `AddinShmemInitLock` when connecting to and initializing its allocation of shared memory, as shown here:

```c
static mystruct *ptr = NULL;
if (!ptr)
{
    bool found;
```
Chapter 36. Extending SQL

LWLockAcquire(AddinShmemInitLock, LW_EXCLUSIVE);
ptr = ShmemInitStruct("my struct name", size, &found);
if (!found)
{
    initialize contents of shmem area;
    acquire any requested LWLocks using:
    ptr->locks = GetNamedLWLockTranche("my tranche name");
}
LWLockRelease(AddinShmemInitLock);

36.9.13. Using C++ for Extensibility

Although the PostgreSQL backend is written in C, it is possible to write extensions in C++ if these guidelines are followed:

- All functions accessed by the backend must present a C interface to the backend; these C functions can then call C++ functions. For example, extern C linkage is required for backend-accessed functions. This is also necessary for any functions that are passed as pointers between the backend and C++ code.
- Free memory using the appropriate deallocation method. For example, most backend memory is allocated using palloc(), so use pfree() to free it. Using C++ delete in such cases will fail.
- Prevent exceptions from propagating into the C code (use a catch-all block at the top level of all extern C functions). This is necessary even if the C++ code does not explicitly throw any exceptions, because events like out-of-memory can still throw exceptions. Any exceptions must be caught and appropriate errors passed back to the C interface. If possible, compile C++ with -fno-exceptions to eliminate exceptions entirely; in such cases, you must check for failures in your C++ code, e.g. check for NULL returned by new().
- If calling backend functions from C++ code, be sure that the C++ call stack contains only plain old data structures (POD). This is necessary because backend errors generate a distant longjmp() that does not properly unroll a C++ call stack with non-POD objects.

In summary, it is best to place C++ code behind a wall of extern C functions that interface to the backend, and avoid exception, memory, and call stack leakage.

36.10. User-defined Aggregates

Aggregate functions in PostgreSQL are defined in terms of state values and state transition functions. That is, an aggregate operates using a state value that is updated as each successive input row is processed. To define a new aggregate function, one selects a data type for the state value, an initial value for the state, and a state transition function. The state transition function takes the previous state value and the aggregate’s input value(s) for the current row, and returns a new state value. A final function can also be specified, in case the desired result of the aggregate is different from the data that needs to be kept in the running state value. The final function takes the ending state value and
returns whatever is wanted as the aggregate result. In principle, the transition and final functions are just ordinary functions that could also be used outside the context of the aggregate. (In practice, it’s often helpful for performance reasons to create specialized transition functions that can only work when called as part of an aggregate.)

Thus, in addition to the argument and result data types seen by a user of the aggregate, there is an internal state-value data type that might be different from both the argument and result types.

If we define an aggregate that does not use a final function, we have an aggregate that computes a running function of the column values from each row. sum is an example of this kind of aggregate. sum starts at zero and always adds the current row’s value to its running total. For example, if we want to make a sum aggregate to work on a data type for complex numbers, we only need the addition function for that data type. The aggregate definition would be:

CREATE AGGREGATE sum (complex)
(
    sfunc = complex_add,
    stype = complex,
    initcond = '(0,0)'
);

which we might use like this:

SELECT sum(a) FROM test_complex;

sum
-----------
(34,53.9)

(Notice that we are relying on function overloading: there is more than one aggregate named sum, but PostgreSQL can figure out which kind of sum applies to a column of type complex.)

The above definition of sum will return zero (the initial state value) if there are no nonnull input values. Perhaps we want to return null in that case instead — the SQL standard expects sum to behave that way. We can do this simply by omitting the initcond phrase, so that the initial state value is null. Ordinarily this would mean that the sfunc would need to check for a null state-value input. But for sum and some other simple aggregates like max and min, it is sufficient to insert the first nonnull input value into the state variable and then start applying the transition function at the second nonnull input value. PostgreSQL will do that automatically if the initial state value is null and the transition function is marked “strict” (i.e., not to be called for null inputs).

Another bit of default behavior for a “strict” transition function is that the previous state value is retained unchanged whenever a null input value is encountered. Thus, null values are ignored. If you need some other behavior for null inputs, do not declare your transition function as strict; instead code it to test for null inputs and do whatever is needed.

avg (average) is a more complex example of an aggregate. It requires two pieces of running state: the sum of the inputs and the count of the number of inputs. The final result is obtained by dividing these quantities. Average is typically implemented by using an array as the state value. For example, the built-in implementation of avg(float8) looks like:

CREATE AGGREGATE avg (float8)
(
    sfunc = float8_accum,
    stype = float8[],
    finalfunc = float8_avg,
    initcond = '{0,0,0}'
);
Chapter 36. Extending SQL

Note: float8_accum requires a three-element array, not just two elements, because it accumulates the sum of squares as well as the sum and count of the inputs. This is so that it can be used for some other aggregates as well as `avg`.

Aggregate function calls in SQL allow `DISTINCT` and `ORDER BY` options that control which rows are fed to the aggregate’s transition function and in what order. These options are implemented behind the scenes and are not the concern of the aggregate’s support functions.

For further details see the CREATE AGGREGATE command.

### 36.10.1. Moving-Aggregate Mode

Aggregate functions can optionally support moving-aggregate mode, which allows substantially faster execution of aggregate functions within windows with moving frame starting points. (See Section 3.5 and Section 4.2.8 for information about use of aggregate functions as window functions.) The basic idea is that in addition to a normal “forward” transition function, the aggregate provides an inverse transition function, which allows rows to be removed from the aggregate’s running state value when they exit the window frame. For example a `sum` aggregate, which uses addition as the forward transition function, would use subtraction as the inverse transition function. Without an inverse transition function, the window function mechanism must recalculate the aggregate from scratch each time the frame starting point moves, resulting in run time proportional to the number of input rows times the average frame length. With an inverse transition function, the run time is only proportional to the number of input rows.

The inverse transition function is passed the current state value and the aggregate input value(s) for the earliest row included in the current state. It must reconstruct what the state value would have been if the given input row had never been aggregated, but only the rows following it. This sometimes requires that the forward transition function keep more state than is needed for plain aggregation mode. Therefore, the moving-aggregate mode uses a completely separate implementation from the plain mode: it has its own state data type, its own forward transition function, and its own final function if needed. These can be the same as the plain mode’s data type and functions, if there is no need for extra state.

As an example, we could extend the `sum` aggregate given above to support moving-aggregate mode like this:

```sql
CREATE AGGREGATE sum (complex)
(
    sfunc = complex_add,
    stype = complex,
    initcond = '(0,0)',
    msfunc = complex_add,
    minvfunc = complex_sub,
    mstype = complex,
    minitcond = '(0,0)'
);
```

The parameters whose names begin with `m` define the moving-aggregate implementation. Except for the inverse transition function `minvfunc`, they correspond to the plain-aggregate parameters without `m`. 

1028
The forward transition function for moving-aggregate mode is not allowed to return null as the new state value. If the inverse transition function returns null, this is taken as an indication that the inverse function cannot reverse the state calculation for this particular input, and so the aggregate calculation will be redone from scratch for the current frame starting position. This convention allows moving-aggregate mode to be used in situations where there are some infrequent cases that are impractical to reverse out of the running state value. The inverse transition function can “punt” on these cases, and yet still come out ahead so long as it can work for most cases. As an example, an aggregate working with floating-point numbers might choose to punt when a \( \text{NaN} \) (not a number) input has to be removed from the running state value.

When writing moving-aggregate support functions, it is important to be sure that the inverse transition function can reconstruct the correct state value exactly. Otherwise there might be user-visible differences in results depending on whether the moving-aggregate mode is used. An example of an aggregate for which adding an inverse transition function seems easy at first, yet where this requirement cannot be met is sum over float4 or float8 inputs. A naive declaration of \( \text{sum(float8)} \) could be

```sql
CREATE AGGREGATE unsafe_sum (float8)
(
    stype = float8,
    sfunc = float8pl,
    mstype = float8,
    msfunc = float8pl,
    minvfunc = float8mi
);
```

This aggregate, however, can give wildly different results than it would have without the inverse transition function. For example, consider

```sql
SELECT unsafe_sum(x) OVER (ORDER BY n ROWS BETWEEN CURRENT ROW AND 1 FOLLOWING)
FROM (VALUES (1, 1.0e20::float8),
           (2, 1.0::float8)) AS v (n,x);
```

This query returns 0 as its second result, rather than the expected answer of 1. The cause is the limited precision of floating-point values: adding 1 to 1e20 results in 1e20 again, and so subtracting 1e20 from that yields 0, not 1. Note that this is a limitation of floating-point arithmetic in general, not a limitation of PostgreSQL.

### 36.10.2. Polymorphic and Variadic Aggregates

Aggregate functions can use polymorphic state transition functions or final functions, so that the same functions can be used to implement multiple aggregates. See Section 36.2.5 for an explanation of polymorphic functions. Going a step further, the aggregate function itself can be specified with polymorphic input type(s) and state type, allowing a single aggregate definition to serve for multiple input data types. Here is an example of a polymorphic aggregate:

```sql
CREATE AGGREGATE array_accum (anyelement)
(
    sfunc = array_append,
    stype = anyarray,
    initcond = '{}'
);
```
Here, the actual state type for any given aggregate call is the array type having the actual input type as elements. The behavior of the aggregate is to concatenate all the inputs into an array of that type. (Note: the built-in aggregate \texttt{array\_agg} provides similar functionality, with better performance than this definition would have.)

Here’s the output using two different actual data types as arguments:

\begin{verbatim}
SELECT attrelid::regclass, array_accum(attname)
FROM pg_attribute
WHERE attnum > 0 AND attrelid = 'pg_tablespace'::regclass
GROUP BY attrelid;
\end{verbatim}

\begin{verbatim}
attrelid | array_accum
----------+---------------------------------------
pg_tablespace | {spcname,spcowner,spcacl,spcoptions}
\end{verbatim}

\begin{verbatim}
(1 row)
\end{verbatim}

\begin{verbatim}
SELECT attrelid::regclass, array_accum(atttypid::regtype)
FROM pg_attribute
WHERE attnum > 0 AND attrelid = 'pg_tablespace'::regclass
GROUP BY attrelid;
\end{verbatim}

\begin{verbatim}
attrelid | array_accum
----------+---------------------------
pg_tablespace | {name,oid,aclitem[],text[]}  
\end{verbatim}

\begin{verbatim}
(1 row)
\end{verbatim}

Ordinarily, an aggregate function with a polymorphic result type has a polymorphic state type, as in the above example. This is necessary because otherwise the final function cannot be declared sensibly: it would need to have a polymorphic result type but no polymorphic argument type, which \texttt{CREATE FUNCTION} will reject on the grounds that the result type cannot be deduced from a call. But sometimes it is inconvenient to use a polymorphic state type. The most common case is where the aggregate support functions are to be written in C and the state type should be declared as \texttt{internal} because there is no SQL-level equivalent for it. To address this case, it is possible to declare the final function as taking extra “dummy” arguments that match the input arguments of the aggregate. Such dummy arguments are always passed as null values since no specific value is available when the final function is called. Their only use is to allow a polymorphic final function’s result type to be connected to the aggregate’s input type(s). For example, the definition of the built-in aggregate \texttt{array\_agg} is equivalent to

\begin{verbatim}
CREATE FUNCTION array_agg_transfn(internal, anynonarray)
    RETURNS internal ...;
CREATE FUNCTION array_agg_finalfn(internal, anynonarray)
    RETURNS anyarray ...;
CREATE AGGREGATE array_agg (anynonarray)
(
    sfunc = array_agg_transfn,
    stype = internal,
    finalfunc = array_agg_finalfn,
    finalfunc_extra
); 
\end{verbatim}
Chapter 36. Extending SQL

Here, the `finalfunc_extra` option specifies that the final function receives, in addition to the state value, extra dummy argument(s) corresponding to the aggregate’s input argument(s). The extra anynonarray argument allows the declaration of `array_agg_finalfn` to be valid.

An aggregate function can be made to accept a varying number of arguments by declaring its last argument as a VARIADIC array, in much the same fashion as for regular functions; see Section 36.4.5. The aggregate’s transition function(s) must have the same array type as their last argument. The transition function(s) typically would also be marked VARIADIC, but this is not strictly required.

**Note:** Variadic aggregates are easily misused in connection with the `ORDER BY` option (see Section 4.2.7), since the parser cannot tell whether the wrong number of actual arguments have been given in such a combination. Keep in mind that everything to the right of `ORDER BY` is a sort key, not an argument to the aggregate. For example, in

```sql
SELECT myaggregate(a ORDER BY a, b, c) FROM ...
```

the parser will see this as a single aggregate function argument and three sort keys. However, the user might have intended

```sql
SELECT myaggregate(a, b, c ORDER BY a) FROM ...
```

If `myaggregate` is variadic, both these calls could be perfectly valid.

For the same reason, it’s wise to think twice before creating aggregate functions with the same names and different numbers of regular arguments.

### 36.10.3. Ordered-Set Aggregates

The aggregates we have been describing so far are “normal” aggregates. PostgreSQL also supports ordered-set aggregates, which differ from normal aggregates in two key ways. First, in addition to ordinary aggregated arguments that are evaluated once per input row, an ordered-set aggregate can have “direct” arguments that are evaluated only once per aggregation operation. Second, the syntax for the ordinary aggregated arguments specifies a sort ordering for them explicitly. An ordered-set aggregate is usually used to implement a computation that depends on a specific row ordering, for instance rank or percentile, so that the sort ordering is a required aspect of any call. For example, the built-in definition of `percentile_disc` is equivalent to:

```sql
CREATE FUNCTION ordered_set_transition(internal, anyelement)
    RETURNS internal ...;
CREATE FUNCTION percentile_disc_final(internal, float8, anyelement)
    RETURNS anyelement ...;

CREATE AGGREGATE percentile_disc (float8 ORDER BY anyelement)
    (
        sfunc = ordered_set_transition,
        stype = internal,
        finalfunc = percentile_disc_final,
        finalfunc_extra
    );
```

This aggregate takes a `float8` direct argument (the percentile fraction) and an aggregated input that can be of any sortable data type. It could be used to obtain a median household income like this:

```sql
SELECT percentile_disc(0.5) WITHIN GROUP (ORDER BY income) FROM households;
```
Here, \(0.5\) is a direct argument; it would make no sense for the percentile fraction to be a value varying across rows.

Unlike the case for normal aggregates, the sorting of input rows for an ordered-set aggregate is \emph{not} done behind the scenes, but is the responsibility of the aggregate's support functions. The typical implementation approach is to keep a reference to a “tuplesort” object in the aggregate’s state value, feed the incoming rows into that object, and then complete the sorting and read out the data in the final function. This design allows the final function to perform special operations such as injecting additional “hypothetical” rows into the data to be sorted. While normal aggregates can often be implemented with support functions written in PL/pgSQL or another PL language, ordered-set aggregates generally have to be written in C, since their state values aren’t definable as any SQL data type. (In the above example, notice that the state value is declared as type \texttt{internal} — this is typical.)

The state transition function for an ordered-set aggregate receives the current state value plus the aggregated input values for each row, and returns the updated state value. This is the same definition as for normal aggregates, but note that the direct arguments (if any) are not provided. The final function receives the last state value, the values of the direct arguments if any, and (if \texttt{finalfunc_extra} is specified) null values corresponding to the aggregated input(s). As with normal aggregates, \texttt{finalfunc_extra} is only really useful if the aggregate is polymorphic; then the extra dummy argument(s) are needed to connect the final function’s result type to the aggregate’s input type(s).

Currently, ordered-set aggregates cannot be used as window functions, and therefore there is no need for them to support moving-aggregate mode.

### 36.10.4. Partial Aggregation

Optionally, an aggregate function can support \emph{partial aggregation}. The idea of partial aggregation is to run the aggregate’s state transition function over different subsets of the input data independently, and then to combine the state values resulting from those subsets to produce the same state value that would have resulted from scanning all the input in a single operation. This mode can be used for parallel aggregation by having different worker processes scan different portions of a table. Each worker produces a partial state value, and at the end those state values are combined to produce a final state value. (In the future this mode might also be used for purposes such as combining aggregations over local and remote tables; but that is not implemented yet.)

To support partial aggregation, the aggregate definition must provide a \emph{combine function}, which takes two values of the aggregate’s state type (representing the results of aggregating over two subsets of the input rows) and produces a new value of the state type, representing what the state would have been after aggregating over the combination of those sets of rows. It is unspecified what the relative order of the input rows from the two sets would have been. This means that it’s usually impossible to define a useful combine function for aggregates that are sensitive to input row order.

As simple examples, \texttt{MAX} and \texttt{MIN} aggregates can be made to support partial aggregation by specifying the combine function as the same greater-of-two or lesser-of-two comparison function that is used as their transition function. \texttt{SUM} aggregates just need an addition function as combine function. (Again, this is the same as their transition function, unless the state value is wider than the input data type.)

The combine function is treated much like a transition function that happens to take a value of the state type, not of the underlying input type, as its second argument. In particular, the rules for dealing with null values and strict functions are similar. Also, if the aggregate definition specifies a non-null
initcond, keep in mind that that will be used not only as the initial state for each partial aggregation run, but also as the initial state for the combine function, which will be called to combine each partial result into that state.

If the aggregate’s state type is declared as internal, it is the combine function’s responsibility that its result is allocated in the correct memory context for aggregate state values. This means in particular that when the first input is NULL it’s invalid to simply return the second input, as that value will be in the wrong context and will not have sufficient lifespan.

When the aggregate’s state type is declared as internal, it is usually also appropriate for the aggregate definition to provide a serialization function and a deserialization function, which allow such a state value to be copied from one process to another. Without these functions, parallel aggregation cannot be performed, and future applications such as local/remote aggregation will probably not work either.

A serialization function must take a single argument of type internal and return a result of type bytea, which represents the state value packaged up into a flat blob of bytes. Conversely, a deserialization function reverses that conversion. It must take two arguments of types bytea and internal, and return a result of type internal. (The second argument is unused and is always zero, but it is required for type-safety reasons.) The result of the deserialization function should simply be allocated in the current memory context, as unlike the combine function’s result, it is not long-lived.

Worth noting also is that for an aggregate to be executed in parallel, the aggregate itself must be marked PARALLEL SAFE. The parallel-safety markings on its support functions are not consulted.

### 36.10.5. Support Functions for Aggregates

A function written in C can detect that it is being called as an aggregate support function by calling `AggCheckCallContext`, for example:

```c
if (AggCheckCallContext(fcinfo, NULL))
```

One reason for checking this is that when it is true for a transition function, the first input must be a temporary state value and can therefore safely be modified in-place rather than allocating a new copy. See `int8inc()` for an example. (This is the only case where it is safe for a function to modify a pass-by-reference input. In particular, final functions for normal aggregates must not modify their inputs in any case, because in some cases they will be re-executed on the same final state value.)

The second argument of AggCheckCallContext can be used to retrieve the memory context in which aggregate state values are being kept. This is useful for transition functions that wish to use “expanded” objects (see Section 36.11.1) as their state values. On first call, the transition function should return an expanded object whose memory context is a child of the aggregate state context, and then keep returning the same expanded object on subsequent calls. See `array_append()` for an example. (array_append() is not the transition function of any built-in aggregate, but it is written to behave efficiently when used as transition function of a custom aggregate.)

Another support routine available to aggregate functions written in C is `AggGetAggref`, which returns the Aggref parse node that defines the aggregate call. This is mainly useful for ordered-set aggregates, which can inspect the substructure of the Aggref node to find out what sort ordering they are supposed to implement. Examples can be found in `orderedsetaggs.c` in the PostgreSQL source code.
36.11. User-defined Types

As described in Section 36.2, PostgreSQL can be extended to support new data types. This section describes how to define new base types, which are data types defined below the level of the SQL language. Creating a new base type requires implementing functions to operate on the type in a low-level language, usually C.

The examples in this section can be found in `complex.sql` and `complex.c` in the `src/tutorial` directory of the source distribution. See the README file in that directory for instructions about running the examples.

A user-defined type must always have input and output functions. These functions determine how the type appears in strings (for input by the user and output to the user) and how the type is organized in memory. The input function takes a null-terminated character string as its argument and returns the internal (in memory) representation of the type. The output function takes the internal representation of the type as argument and returns a null-terminated character string. If we want to do anything more with the type than merely store it, we must provide additional functions to implement whatever operations we’d like to have for the type.

Suppose we want to define a type `complex` that represents complex numbers. A natural way to represent a complex number in memory would be the following C structure:

```c
typedef struct Complex {
    double x;
    double y;
} Complex;
```

We will need to make this a pass-by-reference type, since it’s too large to fit into a single `Datum` value.

As the external string representation of the type, we choose a string of the form `(x,y)`.

The input and output functions are usually not hard to write, especially the output function. But when defining the external string representation of the type, remember that you must eventually write a complete and robust parser for that representation as your input function. For instance:

```c
PG_FUNCTION_INFO_V1(complex_in);

Datum complex_in(PG_FUNCTION_ARGS)
{
    char *str = PG_GETARG_CSTRING(0);
    double x, y;
    Complex *result;

    if (sscanf(str, " ( %lf , %lf )", &x, &y) != 2)
        ereport(ERROR,
            (errcode(ERRCODE_INVALID_TEXT_REPRESENTATION),
             errmsg("invalid input syntax for complex: \"%s\"", str)));

    result = (Complex *) palloc(sizeof(Complex));
    result->x = x;
    result->y = y;
    PG_RETURN_POINTER(result);
}
```

The output function can simply be:
Chapter 36. Extending SQL

PG_FUNCTION_INFO_V1(complex_out);

Datum
complex_out(PG_FUNCTION_ARGS)
{
    Complex  *complex = (Complex *) PG_GETARG_POINTER(0);
    char      *result;

    result = psprintf("(%g,%g)" , complex->x, complex->y);
    PG_RETURN_CSTRING(result);
}

You should be careful to make the input and output functions inverses of each other. If you do not, you will have severe problems when you need to dump your data into a file and then read it back in. This is a particularly common problem when floating-point numbers are involved.

Optionally, a user-defined type can provide binary input and output routines. Binary I/O is normally faster but less portable than textual I/O. As with textual I/O, it is up to you to define exactly what the external binary representation is. Most of the built-in data types try to provide a machine-independent binary representation. For complex, we will piggy-back on the binary I/O converters for type float8:

PG_FUNCTION_INFO_V1(complex_recv);

Datum
complex_recv(PG_FUNCTION_ARGS)
{
    StringInfo buf = (StringInfo) PG_GETARG_POINTER(0);
    Complex   *result;

    result = (Complex *) palloc(sizeof(Complex));
    result->x = pq_getmsgfloat8(buf);
    result->y = pq_getmsgfloat8(buf);
    PG_RETURN_POINTER(result);
}

PG_FUNCTION_INFO_V1(complex_send);

Datum
complex_send(PG_FUNCTION_ARGS)
{
    Complex   *complex = (Complex *) PG_GETARG_POINTER(0);
    StringInfoData buf;

    pq_begintypsend(&buf);
    pq_sendfloat8(&buf, complex->x);
    pq_sendfloat8(&buf, complex->y);
    PG_RETURN_BYTEA_P(pq_endtypsend(&buf));
}

Once we have written the I/O functions and compiled them into a shared library, we can define the complex type in SQL. First we declare it as a shell type:

CREATE TYPE complex;
Chapter 36. Extending SQL

This serves as a placeholder that allows us to reference the type while defining its I/O functions. Now we can define the I/O functions:

```
CREATE FUNCTION complex_in(cstring)
  RETURNS complex
  AS 'filename'
  LANGUAGE C IMMUTABLE STRICT;

CREATE FUNCTION complex_out(complex)
  RETURNS cstring
  AS 'filename'
  LANGUAGE C IMMUTABLE STRICT;

CREATE FUNCTION complex_recv(internal)
  RETURNS complex
  AS 'filename'
  LANGUAGE C IMMUTABLE STRICT;

CREATE FUNCTION complex_send(complex)
  RETURNS bytea
  AS 'filename'
  LANGUAGE C IMMUTABLE STRICT;
```

Finally, we can provide the full definition of the data type:

```
CREATE TYPE complex (
  internallength = 16,
  input = complex_in,
  output = complex_out,
  receive = complex_recv,
  send = complex_send,
  alignment = double
);
```

When you define a new base type, PostgreSQL automatically provides support for arrays of that type. The array type typically has the same name as the base type with the underscore character (_) prepended.

Once the data type exists, we can declare additional functions to provide useful operations on the data type. Operators can then be defined atop the functions, and if needed, operator classes can be created to support indexing of the data type. These additional layers are discussed in following sections.

If the internal representation of the data type is variable-length, the internal representation must follow the standard layout for variable-length data: the first four bytes must be a char[4] field which is never accessed directly (customarily named vl_len_). You must use the SET_VARSIZE() macro to store the total size of the datum (including the length field itself) in this field and VARSIZE() to retrieve it. (These macros exist because the length field may be encoded depending on platform.)

For further details see the description of the CREATE TYPE command.
36.11.1. TOAST Considerations

If the values of your data type vary in size (in internal form), it’s usually desirable to make the data type TOAST-able (see Section 65.2). You should do this even if the values are always too small to be compressed or stored externally, because TOAST can save space on small data too, by reducing header overhead.

To support TOAST storage, the C functions operating on the data type must always be careful to unpack any toasted values they are handed by using `PG_DETOAST_DATUM`. (This detail is customarily hidden by defining type-specific `GETARG_DATATYPE_P` macros.) Then, when running the `CREATE TYPE` command, specify the internal length as `variable` and select some appropriate storage option other than `plain`.

If data alignment is unimportant (either just for a specific function or because the data type specifies byte alignment anyway) then it’s possible to avoid some of the overhead of `PG_DETOAST_DATUM`. You can use `PG_DETOAST_DATUM_PACKED` instead (customarily hidden by defining a `GETARG_DATATYPE_PP` macro) and using the macros `VARSIZE_ANY_EXHDR` and `VARDATA_ANY` to access a potentially-packed datum. Again, the data returned by these macros is not aligned even if the data type definition specifies an alignment. If the alignment is important you must go through the regular `PG_DETOAST_DATUM` interface.

**Note:** Older code frequently declares `vl_len_` as an `int32` field instead of `char[4]`. This is OK as long as the struct definition has other fields that have at least `int32` alignment. But it is dangerous to use such a struct definition when working with a potentially unaligned datum; the compiler may take it as license to assume the datum actually is aligned, leading to core dumps on architectures that are strict about alignment.

Another feature that’s enabled by TOAST support is the possibility of having an expanded in-memory data representation that is more convenient to work with than the format that is stored on disk. The regular or “flat” varlena storage format is ultimately just a blob of bytes; it cannot for example contain pointers, since it may get copied to other locations in memory. For complex data types, the flat format may be quite expensive to work with, so PostgreSQL provides a way to “expand” the flat format into a representation that is more suited to computation, and then pass that format in-memory between functions of the data type.

To use expanded storage, a data type must define an expanded format that follows the rules given in `src/include/utils/expandeddatum.h`, and provide functions to “expand” a flat varlena value into expanded format and “flatten” the expanded format back to the regular varlena representation. Then ensure that all C functions for the data type can accept either representation, possibly by converting one into the other immediately upon receipt. This does not require fixing all existing functions for the data type at once, because the standard `PG_DETOAST_DATUM` macro is defined to convert expanded inputs into regular flat format. Therefore, existing functions that work with the flat varlena format will continue to work, though slightly inefficiently, with expanded inputs; they need not be converted until and unless better performance is important.

C functions that know how to work with an expanded representation typically fall into two categories: those that can only handle expanded format, and those that can handle either expanded or flat varlena inputs. The former are easier to write but may be less efficient overall, because converting a flat input to expanded form for use by a single function may cost more than is saved by operating on the expanded format. When only expanded format need be handled, conversion of flat inputs to expanded form can be hidden inside an argument-fetching macro, so that the function appears no more complex than one working with traditional varlena input. To handle both types of input, write an argument-fetching function that will detoast external, short-header, and compressed varlena inputs, but not expanded inputs. Such a function can be defined as returning a pointer to a union of the
flat varlena format and the expanded format. Callers can use the `VARATT_IS_EXPANDED_HEADER()` macro to determine which format they received.

The TOAST infrastructure not only allows regular varlena values to be distinguished from expanded values, but also distinguishes “read-write” and “read-only” pointers to expanded values. C functions that only need to examine an expanded value, or will only change it in safe and non-semantically-visible ways, need not care which type of pointer they receive. C functions that produce a modified version of an input value are allowed to modify an expanded input value in-place if they receive a read-write pointer, but must not modify the input if they receive a read-only pointer; in that case they have to copy the value first, producing a new value to modify. A C function that has constructed a new expanded value should always return a read-write pointer to it. Also, a C function that is modifying a read-write expanded value in-place should take care to leave the value in a sane state if it fails partway through.

For examples of working with expanded values, see the standard array infrastructure, particularly `src/backend/utils/adt/array_expanded.c`.

### 36.12. User-defined Operators

Every operator is “syntactic sugar” for a call to an underlying function that does the real work; so you must first create the underlying function before you can create the operator. However, an operator is not merely syntactic sugar, because it carries additional information that helps the query planner optimize queries that use the operator. The next section will be devoted to explaining that additional information.

PostgreSQL supports left unary, right unary, and binary operators. Operators can be overloaded; that is, the same operator name can be used for different operators that have different numbers and types of operands. When a query is executed, the system determines the operator to call from the number and types of the provided operands.

Here is an example of creating an operator for adding two complex numbers. We assume we’ve already created the definition of type `complex` (see Section 36.11). First we need a function that does the work, then we can define the operator:

```sql
CREATE FUNCTION complex_add(complex, complex) 
  RETURNS complex 
  AS 'filename', 'complex_add' 
  LANGUAGE C IMMUTABLE STRICT;

CREATE OPERATOR + ( 
  leftarg = complex, 
  rightarg = complex, 
  procedure = complex_add, 
  commutator = + 
);
```

Now we could execute a query like this:

```sql
SELECT (a + b) AS c FROM test_complex;
```

<table>
<thead>
<tr>
<th>c</th>
</tr>
</thead>
<tbody>
<tr>
<td>(5.2,6.05)</td>
</tr>
</tbody>
</table>
We’ve shown how to create a binary operator here. To create unary operators, just omit one of leftarg (for left unary) or rightarg (for right unary). The procedure clause and the argument clauses are the only required items in CREATE OPERATOR. The commutator clause shown in the example is an optional hint to the query optimizer. Further details about commutator and other optimizer hints appear in the next section.

36.13. Operator Optimization Information

A PostgreSQL operator definition can include several optional clauses that tell the system useful things about how the operator behaves. These clauses should be provided whenever appropriate, because they can make for considerable speedups in execution of queries that use the operator. But if you provide them, you must be sure that they are right! Incorrect use of an optimization clause can result in slow queries, subtly wrong output, or other Bad Things. You can always leave out an optimization clause if you are not sure about it; the only consequence is that queries might run slower than they need to.

Additional optimization clauses might be added in future versions of PostgreSQL. The ones described here are all the ones that release 9.6.17 understands.

36.13.1. COMMUTATOR

The COMMUTATOR clause, if provided, names an operator that is the commutator of the operator being defined. We say that operator A is the commutator of operator B if (x A y) equals (y B x) for all possible input values x, y. Notice that B is also the commutator of A. For example, operators < and > for a particular data type are usually each others’ commutators, and operator + is usually commutative with itself. But operator - is usually not commutative with anything.

The left operand type of a commutable operator is the same as the right operand type of its commutator, and vice versa. So the name of the commutator operator is all that PostgreSQL needs to be given to look up the commutator, and that’s all that needs to be provided in the COMMUTATOR clause.

It’s critical to provide commutator information for operators that will be used in indexes and join clauses, because this allows the query optimizer to “flip around” such a clause to the forms needed for different plan types. For example, consider a query with a WHERE clause like tab1.x = tab2.y, where tab1.x and tab2.y are of a user-defined type, and suppose that tab2.y is indexed. The optimizer cannot generate an index scan unless it can determine how to flip the clause around to tab2.y = tab1.x, because the index-scan machinery expects to see the indexed column on the left of the operator it is given. PostgreSQL will not simply assume that this is a valid transformation — the creator of the = operator must specify that it is valid, by marking the operator with commutator information.

When you are defining a self-commutative operator, you just do it. When you are defining a pair of commutative operators, things are a little trickier: how can the first one to be defined refer to the other one, which you haven’t defined yet? There are two solutions to this problem:

- One way is to omit the COMMUTATOR clause in the first operator that you define, and then provide one in the second operator’s definition. Since PostgreSQL knows that commutative operators come in pairs, when it sees the second definition it will automatically go back and fill in the missing COMMUTATOR clause in the first definition.
• The other, more straightforward way is just to include \texttt{COMMUTATOR} clauses in both definitions. When PostgreSQL processes the first definition and realizes that \texttt{COMMUTATOR} refers to a nonexistent operator, the system will make a dummy entry for that operator in the system catalog. This dummy entry will have valid data only for the operator name, left and right operand types, and result type, since that’s all that PostgreSQL can deduce at this point. The first operator’s catalog entry will link to this dummy entry. Later, when you define the second operator, the system updates the dummy entry with the additional information from the second definition. If you try to use the dummy operator before it’s been filled in, you’ll just get an error message.

36.13.2. \texttt{NEGATOR}

The \texttt{NEGATOR} clause, if provided, names an operator that is the negator of the operator being defined. We say that operator A is the negator of operator B if both return Boolean results and (x A y) equals NOT (x B y) for all possible inputs x, y. Notice that B is also the negator of A. For example, < and >= are a negator pair for most data types. An operator can never validly be its own negator.

Unlike commutators, a pair of unary operators could validly be marked as each other’s negators; that would mean (A x) equals NOT (B x) for all x, or the equivalent for right unary operators.

An operator’s negator must have the same left and/or right operand types as the operator to be defined, so just as with \texttt{COMMUTATOR}, only the operator name need be given in the \texttt{NEGATOR} clause.

Providing a negator is very helpful to the query optimizer since it allows expressions like NOT (x = y) to be simplified into x <> y. This comes up more often than you might think, because NOT operations can be inserted as a consequence of other rearrangements.

Pairs of negator operators can be defined using the same methods explained above for commutator pairs.

36.13.3. \texttt{RESTRICT}

The \texttt{RESTRICT} clause, if provided, names a restriction selectivity estimation function for the operator. (Note that this is a function name, not an operator name.) \texttt{RESTRICT} clauses only make sense for binary operators that return boolean. The idea behind a restriction selectivity estimator is to guess what fraction of the rows in a table will satisfy a \texttt{WHERE}-clause condition of the form:

\texttt{column OP constant}

for the current operator and a particular constant value. This assists the optimizer by giving it some idea of how many rows will be eliminated by \texttt{WHERE} clauses that have this form. (What happens if the constant is on the left, you might be wondering? Well, that’s one of the things that \texttt{COMMUTATOR} is for...)

Writing new restriction selectivity estimation functions is far beyond the scope of this chapter, but fortunately you can usually just use one of the system’s standard estimators for many of your own operators. These are the standard restriction estimators:

\begin{verbatim}
eqsel for =
neqsel for <>
scalarltssel for < or <=
scalargtssel for > or >=
\end{verbatim}
It might seem a little odd that these are the categories, but they make sense if you think about it. 
= will typically accept only a small fraction of the rows in a table; <> will typically reject only a 
small fraction. < will accept a fraction that depends on where the given constant falls in the range 
of values for that table column (which, it just so happens, is information collected by ANALYZE and 
made available to the selectivity estimator). <= will accept a slightly larger fraction than < for the 
same comparison constant, but they’re close enough to not be worth distinguishing, especially since 
we’re not likely to do better than a rough guess anyhow. Similar remarks apply to > and >=.

You can frequently get away with using either eqsel or neqsel for operators that have very high or 
very low selectivity, even if they aren’t really equality or inequality. For example, the approximate-
equality geometric operators use eqsel on the assumption that they’ll usually only match a small 
fraction of the entries in a table.

You can use scalarltsel and scalargtssel for comparisons on data types that have 
some sensible means of being converted into numeric scalars for range comparisons. If 
possible, add the data type to those understood by the function convert_to_scalar() in 
src/backend/utils/adt/selfuncs.c. (Eventually, this function should be replaced by 
per-data-type functions identified through a column of the pg_type system catalog; but that hasn’t 
happened yet.) If you do not do this, things will still work, but the optimizer’s estimates won’t be as 
good as they could be.

There are additional selectivity estimation functions designed for geometric operators in 
src/backend/utils/adt/geo_selfuncs.c: areasel, positionsel, and contssel. At this 
writing these are just stubs, but you might want to use them (or even better, improve them) anyway.

### 36.13.4. JOIN

The JOIN clause, if provided, names a join selectivity estimation function for the operator. (Note that 
this is a function name, not an operator name.) JOIN clauses only make sense for binary operators 
that return boolean. The idea behind a join selectivity estimator is to guess what fraction of the rows 
in a pair of tables will satisfy a WHERE-clause condition of the form:

```
table1.column1 OP table2.column2
```

for the current operator. As with the RESTRICT clause, this helps the optimizer very substantially by 
letting it figure out which of several possible join sequences is likely to take the least work.

As before, this chapter will make no attempt to explain how to write a join selectivity estimator 
function, but will just suggest that you use one of the standard estimators if one is applicable:

- eqjoinsel for =
- neqjoinsel for <>
- scalarltjoinsel for < or <=
- scalargtjoinsel for > or >=
- areajoinsel for 2D area-based comparisons
- positionjoinsel for 2D position-based comparisons
- contjoinsel for 2D containment-based comparisons
Chapter 36. Extending SQL

36.13.5. HASHES

The HASHES clause, if present, tells the system that it is permissible to use the hash join method for a join based on this operator. HASHES only makes sense for a binary operator that returns boolean, and in practice the operator must represent equality for some data type or pair of data types.

The assumption underlying hash join is that the join operator can only return true for pairs of left and right values that hash to the same hash code. If two values get put in different hash buckets, the join will never compare them at all, implicitly assuming that the result of the join operator must be false. So it never makes sense to specify HASHES for operators that do not represent some form of equality.

In most cases it is only practical to support hashing for operators that take the same data type on both sides. However, sometimes it is possible to design compatible hash functions for two or more data types; that is, functions that will generate the same hash codes for “equal” values, even though the values have different representations. For example, it’s fairly simple to arrange this property when hashing integers of different widths.

To be marked HASHES, the join operator must appear in a hash index operator family. This is not enforced when you create the operator, since of course the referencing operator family couldn’t exist yet. But attempts to use the operator in hash joins will fail at run time if no such operator family exists. The system needs the operator family to find the data-type-specific hash function(s) for the operator’s input data type(s). Of course, you must also create suitable hash functions before you can create the operator family.

Care should be exercised when preparing a hash function, because there are machine-dependent ways in which it might fail to do the right thing. For example, if your data type is a structure in which there might be uninteresting pad bits, you cannot simply pass the whole structure to hash_any. (Unless you write your other operators and functions to ensure that the unused bits are always zero, which is the recommended strategy.) Another example is that on machines that meet the IEEE floating-point standard, negative zero and positive zero are different values (different bit patterns) but they are defined to compare equal. If a float value might contain negative zero then extra steps are needed to ensure it generates the same hash value as positive zero.

A hash-joinable operator must have a commutator (itself if the two operand data types are the same, or a related equality operator if they are different) that appears in the same operator family. If this is not the case, planner errors might occur when the operator is used. Also, it is a good idea (but not strictly required) for a hash operator family that supports multiple data types to provide equality operators for every combination of the data types; this allows better optimization.

Note: The function underlying a hash-joinable operator must be marked immutable or stable. If it is volatile, the system will never attempt to use the operator for a hash join.

Note: If a hash-joinable operator has an underlying function that is marked strict, the function must also be complete: that is, it should return true or false, never null, for any two nonnull inputs. If this rule is not followed, hash-optimization of IN operations might generate wrong results. (Specifically, IN might return false where the correct answer according to the standard would be null; or it might yield an error complaining that it wasn’t prepared for a null result.)

36.13.6. MERGES

The MERGES clause, if present, tells the system that it is permissible to use the merge-join method for
a join based on this operator. MERGES only makes sense for a binary operator that returns boolean, and in practice the operator must represent equality for some data type or pair of data types.

Merge join is based on the idea of sorting the left- and right-hand tables into order and then scanning them in parallel. So, both data types must be capable of being fully ordered, and the join operator must be one that can only succeed for pairs of values that fall at the “same place” in the sort order. In practice this means that the join operator must behave like equality. But it is possible to merge-join two distinct data types so long as they are logically compatible. For example, the smallint-versus-integer equality operator is merge-joinable. We only need sorting operators that will bring both data types into a logically compatible sequence.

To be marked MERGES, the join operator must appear as an equality member of a btree index operator family. This is not enforced when you create the operator, since of course the referencing operator family couldn’t exist yet. But the operator will not actually be used for merge joins unless a matching operator family can be found. The MERGES flag thus acts as a hint to the planner that it’s worth looking for a matching operator family.

A merge-joinable operator must have a commutator (itself if the two operand data types are the same, or a related equality operator if they are different) that appears in the same operator family. If this is not the case, planner errors might occur when the operator is used. Also, it is a good idea (but not strictly required) for a btree operator family that supports multiple data types to provide equality operators for every combination of the data types; this allows better optimization.

Note: The function underlying a merge-joinable operator must be marked immutable or stable. If it is volatile, the system will never attempt to use the operator for a merge join.

36.14. Interfacing Extensions To Indexes

The procedures described thus far let you define new types, new functions, and new operators. However, we cannot yet define an index on a column of a new data type. To do this, we must define an operator class for the new data type. Later in this section, we will illustrate this concept in an example: a new operator class for the B-tree index method that stores and sorts complex numbers in ascending absolute value order.

Operator classes can be grouped into operator families to show the relationships between semantically compatible classes. When only a single data type is involved, an operator class is sufficient, so we’ll focus on that case first and then return to operator families.

36.14.1. Index Methods and Operator Classes

The pg_am table contains one row for every index method (internally known as access method). Support for regular access to tables is built into PostgreSQL, but all index methods are described in pg_am. It is possible to add a new index access method by writing the necessary code and then creating a row in pg_am — but that is beyond the scope of this chapter (see Chapter 59).

The routines for an index method do not directly know anything about the data types that the index method will operate on. Instead, an operator class identifies the set of operations that the index method needs to use to work with a particular data type. Operator classes are so called because one thing they specify is the set of WHERE-clause operators that can be used with an index (i.e., can be converted into an index-scan qualification). An operator class can also specify some support procedures that
are needed by the internal operations of the index method, but do not directly correspond to any
WHERE-clause operator that can be used with the index.

It is possible to define multiple operator classes for the same data type and index method. By doing
this, multiple sets of indexing semantics can be defined for a single data type. For example, a B-tree
index requires a sort ordering to be defined for each data type it works on. It might be useful for a
complex-number data type to have one B-tree operator class that sorts the data by complex absolute
value, another that sorts by real part, and so on. Typically, one of the operator classes will be deemed
most commonly useful and will be marked as the default operator class for that data type and index
method.

The same operator class name can be used for several different index methods (for example, both
B-tree and hash index methods have operator classes named int4_ops), but each such class is an
independent entity and must be defined separately.

### 36.14.2. Index Method Strategies

The operators associated with an operator class are identified by “strategy numbers”, which serve to
identify the semantics of each operator within the context of its operator class. For example, B-trees
impose a strict ordering on keys, lesser to greater, and so operators like “less than” and “greater than
or equal to” are interesting with respect to a B-tree. Because PostgreSQL allows the user to define
operators, PostgreSQL cannot look at the name of an operator (e.g., `<` or `>=`) and tell what kind
of comparison it is. Instead, the index method defines a set of “strategies”, which can be thought
of as generalized operators. Each operator class specifies which actual operator corresponds to each
strategy for a particular data type and interpretation of the index semantics.

The B-tree index method defines five strategies, shown in Table 36-2.

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>less than</td>
<td>1</td>
</tr>
<tr>
<td>less than or equal</td>
<td>2</td>
</tr>
<tr>
<td>equal</td>
<td>3</td>
</tr>
<tr>
<td>greater than or equal</td>
<td>4</td>
</tr>
<tr>
<td>greater than</td>
<td>5</td>
</tr>
</tbody>
</table>

Hash indexes support only equality comparisons, and so they use only one strategy, shown in Table
36-3.

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>equal</td>
<td>1</td>
</tr>
</tbody>
</table>

GiST indexes are more flexible: they do not have a fixed set of strategies at all. Instead, the “consis-
tency” support routine of each particular GiST operator class interprets the strategy numbers however
it likes. As an example, several of the built-in GiST index operator classes index two-dimensional
geometric objects, providing the “R-tree” strategies shown in Table 36-4. Four of these are true two-
dimensional tests (overlaps, same, contains, contained by); four of them consider only the X direction;
and the other four provide the same tests in the Y direction.

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>equal</td>
<td>1</td>
</tr>
</tbody>
</table>
Table 36-4. GiST Two-Dimensional “R-tree” Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>strictly left of</td>
<td>1</td>
</tr>
<tr>
<td>does not extend to right of</td>
<td>2</td>
</tr>
<tr>
<td>overlaps</td>
<td>3</td>
</tr>
<tr>
<td>does not extend to left of</td>
<td>4</td>
</tr>
<tr>
<td>strictly right of</td>
<td>5</td>
</tr>
<tr>
<td>same</td>
<td>6</td>
</tr>
<tr>
<td>contains</td>
<td>7</td>
</tr>
<tr>
<td>contained by</td>
<td>8</td>
</tr>
<tr>
<td>does not extend above</td>
<td>9</td>
</tr>
<tr>
<td>strictly below</td>
<td>10</td>
</tr>
<tr>
<td>strictly above</td>
<td>11</td>
</tr>
<tr>
<td>does not extend below</td>
<td>12</td>
</tr>
</tbody>
</table>

SP-GiST indexes are similar to GiST indexes in flexibility: they don’t have a fixed set of strategies. Instead the support routines of each operator class interpret the strategy numbers according to the operator class’s definition. As an example, the strategy numbers used by the built-in operator classes for points are shown in Table 36-5.

Table 36-5. SP-GiST Point Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>strictly left of</td>
<td>1</td>
</tr>
<tr>
<td>strictly right of</td>
<td>5</td>
</tr>
<tr>
<td>same</td>
<td>6</td>
</tr>
<tr>
<td>contained by</td>
<td>8</td>
</tr>
<tr>
<td>strictly below</td>
<td>10</td>
</tr>
<tr>
<td>strictly above</td>
<td>11</td>
</tr>
</tbody>
</table>

GIN indexes are similar to GiST and SP-GiST indexes, in that they don’t have a fixed set of strategies either. Instead the support routines of each operator class interpret the strategy numbers according to the operator class’s definition. As an example, the strategy numbers used by the built-in operator classes for arrays are shown in Table 36-6.

Table 36-6. GIN Array Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>overlap</td>
<td>1</td>
</tr>
<tr>
<td>contains</td>
<td>2</td>
</tr>
<tr>
<td>is contained by</td>
<td>3</td>
</tr>
<tr>
<td>equal</td>
<td>4</td>
</tr>
</tbody>
</table>

BRIN indexes are similar to GiST, SP-GiST and GIN indexes in that they don’t have a fixed set of strategies either. Instead the support routines of each operator class interpret the strategy numbers according to the operator class’s definition. As an example, the strategy numbers used by the built-in Minmax operator classes are shown in Table 36-7.
Chapter 36. Extending SQL

Table 36-7. BRIN Minmax Strategies

<table>
<thead>
<tr>
<th>Operation</th>
<th>Strategy Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>less than</td>
<td>1</td>
</tr>
<tr>
<td>less than or equal</td>
<td>2</td>
</tr>
<tr>
<td>equal</td>
<td>3</td>
</tr>
<tr>
<td>greater than or equal</td>
<td>4</td>
</tr>
<tr>
<td>greater than</td>
<td>5</td>
</tr>
</tbody>
</table>

Notice that all the operators listed above return Boolean values. In practice, all operators defined as index method search operators must return type `boolean`, since they must appear at the top level of a `WHERE` clause to be used with an index. (Some index access methods also support ordering operators, which typically don’t return Boolean values; that feature is discussed in Section 36.14.7.)

### 36.14.3. Index Method Support Routines

Strategies aren’t usually enough information for the system to figure out how to use an index. In practice, the index methods require additional support routines in order to work. For example, the B-tree index method must be able to compare two keys and determine whether one is greater than, equal to, or less than the other. Similarly, the hash index method must be able to compute hash codes for key values. These operations do not correspond to operators used in qualifications in SQL commands; they are administrative routines used by the index methods, internally.

Just as with strategies, the operator class identifies which specific functions should play each of these roles for a given data type and semantic interpretation. The index method defines the set of functions it needs, and the operator class identifies the correct functions to use by assigning them to the “support function numbers” specified by the index method.

B-trees require a single support function, and allow a second one to be supplied at the operator class author’s option, as shown in Table 36-8.

Table 36-8. B-tree Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>Compare two keys and return an integer less than zero, zero, or greater than zero, indicating whether the first key is less than, equal to, or greater than the second</td>
<td>1</td>
</tr>
<tr>
<td>Return the addresses of C-callable sort support function(s), as documented in <code>utils/sortsupport.h</code> (optional)</td>
<td>2</td>
</tr>
</tbody>
</table>

Hash indexes require one support function, shown in Table 36-9.

Table 36-9. Hash Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>Compute the hash value for a key</td>
<td>1</td>
</tr>
</tbody>
</table>

GiST indexes have nine support functions, two of which are optional, as shown in Table 36-10. (For more information see Chapter 61.)
Table 36-10. GiST Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>consistent</td>
<td>determine whether key satisfies the query qualifier</td>
<td>1</td>
</tr>
<tr>
<td>union</td>
<td>compute union of a set of keys</td>
<td>2</td>
</tr>
<tr>
<td>compress</td>
<td>compute a compressed representation of a key or value to be indexed</td>
<td>3</td>
</tr>
<tr>
<td>decompress</td>
<td>compute a decompressed representation of a compressed key</td>
<td>4</td>
</tr>
<tr>
<td>penalty</td>
<td>compute penalty for inserting new key into subtree with given subtree’s key</td>
<td>5</td>
</tr>
<tr>
<td>picksplit</td>
<td>determine which entries of a page are to be moved to the new page and compute the union keys for resulting pages</td>
<td>6</td>
</tr>
<tr>
<td>equal</td>
<td>compare two keys and return true if they are equal</td>
<td>7</td>
</tr>
<tr>
<td>distance</td>
<td>determine distance from key to query value (optional)</td>
<td>8</td>
</tr>
<tr>
<td>fetch</td>
<td>compute original representation of a compressed key for index-only scans (optional)</td>
<td>9</td>
</tr>
</tbody>
</table>

SP-GiST indexes require five support functions, as shown in Table 36-11. (For more information see Chapter 62.)

Table 36-11. SP-GiST Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>config</td>
<td>provide basic information about the operator class</td>
<td>1</td>
</tr>
<tr>
<td>choose</td>
<td>determine how to insert a new value into an inner tuple</td>
<td>2</td>
</tr>
<tr>
<td>picksplit</td>
<td>determine how to partition a set of values</td>
<td>3</td>
</tr>
<tr>
<td>inner_consistent</td>
<td>determine which sub-partitions need to be searched for a query</td>
<td>4</td>
</tr>
<tr>
<td>leaf_consistent</td>
<td>determine whether key satisfies the query qualifier</td>
<td>5</td>
</tr>
</tbody>
</table>

GIN indexes have six support functions, three of which are optional, as shown in Table 36-12. (For more information see Chapter 63.)

Table 36-12. GIN Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>config</td>
<td>provide basic information about the operator class</td>
<td>1</td>
</tr>
<tr>
<td>choose</td>
<td>determine how to insert a new value into an inner tuple</td>
<td>2</td>
</tr>
<tr>
<td>inner_consistent</td>
<td>determine whether key satisfies the query qualifier</td>
<td>4</td>
</tr>
<tr>
<td>leaf_consistent</td>
<td>determine whether key satisfies the query qualifier</td>
<td>5</td>
</tr>
</tbody>
</table>
### Table 36-13. BRIN Support Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
<th>Support Number</th>
</tr>
</thead>
<tbody>
<tr>
<td>opcInfo</td>
<td>return internal information describing the indexed columns’ summary data</td>
<td>1</td>
</tr>
<tr>
<td>add_value</td>
<td>add a new value to an existing summary index tuple</td>
<td>2</td>
</tr>
<tr>
<td>consistent</td>
<td>determine whether value matches query condition</td>
<td>3</td>
</tr>
<tr>
<td>union</td>
<td>compute union of two summary tuples</td>
<td>4</td>
</tr>
</tbody>
</table>

BRIN indexes have four basic support functions, as shown in Table 36-13; those basic functions may require additional support functions to be provided. (For more information see Section 64.3.)

Unlike search operators, support functions return whichever data type the particular index method expects; for example in the case of the comparison function for B-trees, a signed integer. The number and types of the arguments to each support function are likewise dependent on the index method. For B-tree and hash the comparison and hashing support functions take the same input data types as do the operators included in the operator class, but this is not the case for most GiST, SP-GiST, GIN, and BRIN support functions.
36.14.4. An Example

Now that we have seen the ideas, here is the promised example of creating a new operator class. (You can find a working copy of this example in src/tutorial/complex.c and src/tutorial/complex.sql in the source distribution.) The operator class encapsulates operators that sort complex numbers in absolute value order, so we choose the name complex_abs_ops. First, we need a set of operators. The procedure for defining operators was discussed in Section 36.12. For an operator class on B-trees, the operators we require are:

- absolute-value less-than (strategy 1)
- absolute-value less-than-or-equal (strategy 2)
- absolute-value equal (strategy 3)
- absolute-value greater-than-or-equal (strategy 4)
- absolute-value greater-than (strategy 5)

The least error-prone way to define a related set of comparison operators is to write the B-tree comparison support function first, and then write the other functions as one-line wrappers around the support function. This reduces the odds of getting inconsistent results for corner cases. Following this approach, we first write:

```c
#define Mag(c) ((c)->x *(c)->x + (c)->y*(c)->y)
static int complex_abs_cmp_internal(Complex *a, Complex *b) {
    double amag = Mag(a),
            bmag = Mag(b);

    if (amag < bmag)
        return -1;
    if (amag > bmag)
        return 1;
    return 0;
}
```

Now the less-than function looks like:

```c
PG_FUNCTION_INFO_V1(complex_abs_lt);
Datum complex_abs_lt(PG_FUNCTION_ARGS) {
    Complex *a = (Complex *) PG_GETARG_POINTER(0);
    Complex *b = (Complex *) PG_GETARG_POINTER(1);
    PG_RETURN_BOOL(complex_abs_cmp_internal(a, b) < 0);
}
```

The other four functions differ only in how they compare the internal function’s result to zero. Next we declare the functions and the operators based on the functions to SQL:

```sql
CREATE FUNCTION complex_abs_lt(complex, complex) RETURNS bool
    AS 'filename', 'complex_abs_lt'
    LANGUAGE C IMMUTABLE STRICT;
```
CREATE OPERATOR < { 
    leftarg = complex, rightarg = complex, procedure = complex_abs_lt,
    commutator = >, negator = >=,
    restrict = scalarltssel, join = scalarltjoinsel
};

It is important to specify the correct commutator and negator operators, as well as suitable restriction and join selectivity functions, otherwise the optimizer will be unable to make effective use of the index. Note that the less-than, equal, and greater-than cases should use different selectivity functions.

Other things worth noting are happening here:

- There can only be one operator named, say, = and taking type complex for both operands. In this case we don’t have any other operator = for complex, but if we were building a practical data type we’d probably want = to be the ordinary equality operation for complex numbers (and not the equality of the absolute values). In that case, we’d need to use some other operator name for complex_abs_eq.

- Although PostgreSQL can cope with functions having the same SQL name as long as they have different argument data types, C can only cope with one global function having a given name. So we shouldn’t name the C function something simple like abs_eq. Usually it’s a good practice to include the data type name in the C function name, so as not to conflict with functions for other data types.

- We could have made the SQL name of the function abs_eq, relying on PostgreSQL to distinguish it by argument data types from any other SQL function of the same name. To keep the example simple, we make the function have the same names at the C level and SQL level.

The next step is the registration of the support routine required by B-trees. The example C code that implements this is in the same file that contains the operator functions. This is how we declare the function:

```
CREATE FUNCTION complex_abs_cmp(complex, complex)
    RETURNS integer
    AS 'filename'
    LANGUAGE C IMMUTABLE STRICT;
```

Now that we have the required operators and support routine, we can finally create the operator class:

```
CREATE OPERATOR CLASS complex_abs_ops
    DEFAULT FOR TYPE complex USING btree AS 
    OPERATOR 1  <    ,
    OPERATOR 2  <=  ,
    OPERATOR 3  =   ,
    OPERATOR 4  >=  ,
    OPERATOR 5  >   ,
    FUNCTION 1  complex_abs_cmp(complex, complex);
```

And we’re done! It should now be possible to create and use B-tree indexes on complex columns.

We could have written the operator entries more verbosely, as in:

```
OPERATOR 1  < (complex, complex) ,
```
but there is no need to do so when the operators take the same data type we are defining the operator class for.

The above example assumes that you want to make this new operator class the default B-tree operator class for the complex data type. If you don’t, just leave out the word DEFAULT.

### 36.14.5. Operator Classes and Operator Families

So far we have implicitly assumed that an operator class deals with only one data type. While there certainly can be only one data type in a particular index column, it is often useful to index operations that compare an indexed column to a value of a different data type. Also, if there is use for a cross-data-type operator in connection with an operator class, it is often the case that the other data type has a related operator class of its own. It is helpful to make the connections between related classes explicit, because this can aid the planner in optimizing SQL queries (particularly for B-tree operator classes, since the planner contains a great deal of knowledge about how to work with them).

To handle these needs, PostgreSQL uses the concept of an operator family. An operator family contains one or more operator classes, and can also contain indexable operators and corresponding support functions that belong to the family as a whole but not to any single class within the family. We say that such operators and functions are “loose” within the family, as opposed to being bound into a specific class. Typically each operator class contains single-data-type operators while cross-data-type operators are loose in the family.

All the operators and functions in an operator family must have compatible semantics, where the compatibility requirements are set by the index method. You might therefore wonder why bother to single out particular subsets of the family as operator classes; and indeed for many purposes the class divisions are irrelevant and the family is the only interesting grouping. The reason for defining operator classes is that they specify how much of the family is needed to support any particular index. If there is an index using an operator class, then that operator class cannot be dropped without dropping the index — but other parts of the operator family, namely other operator classes and loose operators, could be dropped. Thus, an operator class should be specified to contain the minimum set of operators and functions that are reasonably needed to work with an index on a specific data type, and then related but non-essential operators can be added as loose members of the operator family.

As an example, PostgreSQL has a built-in B-tree operator family integer_ops, which includes operator classes int8_ops, int4_ops, and int2_ops for indexes on bigint (int8), integer (int4), and smallint (int2) columns respectively. The family also contains cross-data-type comparison operators allowing any two of these types to be compared, so that an index on one of these types can be searched using a comparison value of another type. The family could be duplicated by these definitions:

```sql
CREATE OPERATOR FAMILY integer_ops USING btree;

CREATE OPERATOR CLASS int8_ops
DEFAULT FOR TYPE int8 USING btree FAMILY integer_ops AS
  -- standard int8 comparisons
  OPERATOR 1 < ,
  OPERATOR 2 <= ,
  OPERATOR 3 = ,
  OPERATOR 4 >= ,
  OPERATOR 5 > ,
  FUNCTION 1 btint8cmp(int8, int8) ,
  FUNCTION 2 btint8sortsupport(internal) ;

CREATE OPERATOR CLASS int4_ops
```

1051
Chapter 36. Extending SQL

DEFAULT FOR TYPE int4 USING btree FAMILY integer_ops AS
-- standard int4 comparisons
OPERATOR 1 <,
OPERATOR 2 <=,
OPERATOR 3 =,
OPERATOR 4 >=,
OPERATOR 5 >,
FUNCTION 1 btint4cmp(int4, int4),
FUNCTION 2 btint4sortsupport(internal);

CREATE OPERATOR CLASS int2_ops
DEFAULT FOR TYPE int2 USING btree FAMILY integer_ops AS
-- standard int2 comparisons
OPERATOR 1 <,
OPERATOR 2 <=,
OPERATOR 3 =,
OPERATOR 4 >=,
OPERATOR 5 >,
FUNCTION 1 btint2cmp(int2, int2),
FUNCTION 2 btint2sortsupport(internal);

ALTER OPERATOR FAMILY integer_ops USING btree ADD
-- cross-type comparisons int8 vs int2
OPERATOR 1 < (int8, int2),
OPERATOR 2 <= (int8, int2),
OPERATOR 3 = (int8, int2),
OPERATOR 4 >= (int8, int2),
OPERATOR 5 > (int8, int2),
FUNCTION 1 btint82cmp(int8, int2),

-- cross-type comparisons int8 vs int4
OPERATOR 1 < (int8, int4),
OPERATOR 2 <= (int8, int4),
OPERATOR 3 = (int8, int4),
OPERATOR 4 >= (int8, int4),
OPERATOR 5 > (int8, int4),
FUNCTION 1 btint84cmp(int8, int4),

-- cross-type comparisons int4 vs int2
OPERATOR 1 < (int4, int2),
OPERATOR 2 <= (int4, int2),
OPERATOR 3 = (int4, int2),
OPERATOR 4 >= (int4, int2),
OPERATOR 5 > (int4, int2),
FUNCTION 1 btint42cmp(int4, int2),

-- cross-type comparisons int4 vs int8
OPERATOR 1 < (int4, int8),
OPERATOR 2 <= (int4, int8),
OPERATOR 3 = (int4, int8),
OPERATOR 4 >= (int4, int8),
OPERATOR 5 > (int4, int8),
FUNCTION 1 btint48cmp(int4, int8),

-- cross-type comparisons int2 vs int8
OPERATOR 1 < (int2, int8),
OPERATOR 2 <= (int2, int8),
OPERATOR 3 = (int2, int8) ,
OPERATOR 4 >= (int2, int8) ,
OPERATOR 5 > (int2, int8) ,
FUNCTION 1 btint28cmp(int2, int8) ,

-- cross-type comparisons int2 vs int4
OPERATOR 1 < (int2, int4) ,
OPERATOR 2 <= (int2, int4) ,
OPERATOR 3 = (int2, int4) ,
OPERATOR 4 >= (int2, int4) ,
OPERATOR 5 > (int2, int4) ,
FUNCTION 1 btint24cmp(int2, int4) ;

Notice that this definition “overloads” the operator strategy and support function numbers: each number occurs multiple times within the family. This is allowed so long as each instance of a particular number has distinct input data types. The instances that have both input types equal to an operator class’s input type are the primary operators and support functions for that operator class, and in most cases should be declared as part of the operator class rather than as loose members of the family.

In a B-tree operator family, all the operators in the family must sort compatibly, meaning that the transitive laws hold across all the data types supported by the family: “if A = B and B = C, then A = C”, and “if A < B and B < C, then A < C”. Moreover, implicit or binary coercion casts between types represented in the operator family must not change the associated sort ordering. For each operator in the family there must be a support function having the same two input data types as the operator. It is recommended that a family be complete, i.e., for each combination of data types, all operators are included. Each operator class should include just the non-cross-type operators and support function for its data type.

To build a multiple-data-type hash operator family, compatible hash support functions must be created for each data type supported by the family. Here compatibility means that the functions are guaranteed to return the same hash code for any two values that are considered equal by the family’s equality operators, even when the values are of different types. This is usually difficult to accomplish when the types have different physical representations, but it can be done in some cases. Furthermore, casting a value from one data type represented in the operator family to another data type also represented in the operator family via an implicit or binary coercion cast must not change the computed hash value. Notice that there is only one support function per data type, not one per equality operator. It is recommended that a family be complete, i.e., provide an equality operator for each combination of data types. Each operator class should include just the non-cross-type equality operator and the support function for its data type.

GiST, SP-GiST, and GIN indexes do not have any explicit notion of cross-data-type operations. The set of operators supported is just whatever the primary support functions for a given operator class can handle.

In BRIN, the requirements depends on the framework that provides the operator classes. For operator classes based on minmax, the behavior required is the same as for B-tree operator families: all the operators in the family must sort compatibly, and casts must not change the associated sort ordering.

Note: Prior to PostgreSQL 8.3, there was no concept of operator families, and so any cross-data-type operators intended to be used with an index had to be bound directly into the index’s operator class. While this approach still works, it is deprecated because it makes an index’s dependencies too broad, and because the planner can handle cross-data-type comparisons more effectively when both data types have operators in the same operator family.
36.14.6. System Dependencies on Operator Classes

PostgreSQL uses operator classes to infer the properties of operators in more ways than just whether they can be used with indexes. Therefore, you might want to create operator classes even if you have no intention of indexing any columns of your data type.

In particular, there are SQL features such as ORDER BY and DISTINCT that require comparison and sorting of values. To implement these features on a user-defined data type, PostgreSQL looks for the default B-tree operator class for the data type. The “equals” member of this operator class defines the system’s notion of equality of values for GROUP BY and DISTINCT, and the sort ordering imposed by the operator class defines the default ORDER BY ordering.

Comparison of arrays of user-defined types also relies on the semantics defined by the default B-tree operator class.

If there is no default B-tree operator class for a data type, the system will look for a default hash operator class. But since that kind of operator class only provides equality, in practice it is only enough to support array equality.

When there is no default operator class for a data type, you will get errors like “could not identify an ordering operator” if you try to use these SQL features with the data type.

Note: In PostgreSQL versions before 7.4, sorting and grouping operations would implicitly use operators named =, <, and >. The new behavior of relying on default operator classes avoids having to make any assumption about the behavior of operators with particular names.

Another important point is that an operator that appears in a hash operator family is a candidate for hash joins, hash aggregation, and related optimizations. The hash operator family is essential here since it identifies the hash function(s) to use.

36.14.7. Ordering Operators

Some index access methods (currently, only GiST) support the concept of ordering operators. What we have been discussing so far are search operators. A search operator is one for which the index can be searched to find all rows satisfying WHERE indexed_column operator constant. Note that nothing is promised about the order in which the matching rows will be returned. In contrast, an ordering operator does not restrict the set of rows that can be returned, but instead determines their order. An ordering operator is one for which the index can be scanned to return rows in the order represented by ORDER BY indexed_column operator constant. The reason for defining ordering operators that way is that it supports nearest-neighbor searches, if the operator is one that measures distance. For example, a query like

```
SELECT * FROM places ORDER BY location <-> point ' (101,456)' LIMIT 10;
```

finds the ten places closest to a given target point. A GiST index on the location column can do this efficiently because <-> is an ordering operator.

While search operators have to return Boolean results, ordering operators usually return some other type, such as float or numeric for distances. This type is normally not the same as the data type being indexed. To avoid hard-wiring assumptions about the behavior of different data types, the definition of an ordering operator is required to name a B-tree operator family that specifies the sort ordering of the result data type. As was stated in the previous section, B-tree operator families define PostgreSQL’s notion of ordering, so this is a natural representation. Since the point <-> operator returns float8, it could be specified in an operator class creation command like this:
where float_ops is the built-in operator family that includes operations on float8. This declaration states that the index is able to return rows in order of increasing values of the <-> operator.

### 36.14.8. Special Features of Operator Classes

There are two special features of operator classes that we have not discussed yet, mainly because they are not useful with the most commonly used index methods.

Normally, declaring an operator as a member of an operator class (or family) means that the index method can retrieve exactly the set of rows that satisfy a WHERE condition using the operator. For example:

```sql
SELECT * FROM table WHERE integer_column < 4;
```

can be satisfied exactly by a B-tree index on the integer column. But there are cases where an index is useful as an inexact guide to the matching rows. For example, if a GiST index stores only bounding boxes for geometric objects, then it cannot exactly satisfy a WHERE condition that tests overlap between nonrectangular objects such as polygons. Yet we could use the index to find objects whose bounding box overlaps the bounding box of the target object, and then do the exact overlap test only on the objects found by the index. If this scenario applies, the index is said to be “lossy” for the operator. Lossy index searches are implemented by having the index method return a recheck flag when a row might or might not really satisfy the query condition. The core system will then test the original query condition on the retrieved row to see whether it should be returned as a valid match. This approach works if the index is guaranteed to return all the required rows, plus perhaps some additional rows, which can be eliminated by performing the original operator invocation. The index methods that support lossy searches (currently, GiST, SP-GiST and GIN) allow the support functions of individual operator classes to set the recheck flag, and so this is essentially an operator-class feature.

Consider again the situation where we are storing in the index only the bounding box of a complex object such as a polygon. In this case there’s not much value in storing the whole polygon in the index entry — we might as well store just a simpler object of type box. This situation is expressed by the STORAGE option in CREATE OPERATOR CLASS: we’d write something like:

```sql
CREATE OPERATOR CLASS polygon_ops
    DEFAULT FOR TYPE polygon USING gist AS
    ...
    STORAGE box;
```

At present, only the GiST, GIN and BRIN index methods support a STORAGE type that’s different from the column data type. The GiST `compress` and `decompress` support routines must deal with data-type conversion when STORAGE is used. In GIN, the STORAGE type identifies the type of the “key” values, which normally is different from the type of the indexed column — for example, an operator class for integer-array columns might have keys that are just integers. The GIN `extractValue` and `extractQuery` support routines are responsible for extracting keys from indexed values. BRIN is similar to GIN: the STORAGE type identifies the type of the stored summary values, and operator classes’ support procedures are responsible for interpreting the summary values correctly.
36.15. Packaging Related Objects into an Extension

A useful extension to PostgreSQL typically includes multiple SQL objects; for example, a new data type will require new functions, new operators, and probably new index operator classes. It is helpful to collect all these objects into a single package to simplify database management. PostgreSQL calls such a package an extension. To define an extension, you need at least a script file that contains the SQL commands to create the extension’s objects, and a control file that specifies a few basic properties of the extension itself. If the extension includes C code, there will typically also be a shared library file into which the C code has been built. Once you have these files, a simple CREATE EXTENSION command loads the objects into your database.

The main advantage of using an extension, rather than just running the SQL script to load a bunch of “loose” objects into your database, is that PostgreSQL will then understand that the objects of the extension go together. You can drop all the objects with a single DROP EXTENSION command (no need to maintain a separate “uninstall” script). Even more useful, pg_dump knows that it should not dump the individual member objects of the extension — it will just include a CREATE EXTENSION command in dumps, instead. This vastly simplifies migration to a new version of the extension that might contain more or different objects than the old version. Note however that you must have the extension’s control, script, and other files available when loading such a dump into a new database.

PostgreSQL will not let you drop an individual object contained in an extension, except by dropping the whole extension. Also, while you can change the definition of an extension member object (for example, via CREATE OR REPLACE FUNCTION for a function), bear in mind that the modified definition will not be dumped by pg_dump. Such a change is usually only sensible if you concurrently make the same change in the extension’s script file. (But there are special provisions for tables containing configuration data; see Section 36.15.4.) In production situations, it’s generally better to create an extension update script to perform changes to extension member objects.

The extension script may set privileges on objects that are part of the extension via GRANT and REVOKE statements. The final set of privileges for each object (if any are set) will be stored in the pg_init_privs system catalog. When pg_dump is used, the CREATE EXTENSION command will be included in the dump, followed by the set of GRANT and REVOKE statements necessary to set the privileges on the objects to what they were at the time the dump was taken.

PostgreSQL does not currently support extension scripts issuing CREATE POLICY or SECURITY LABEL statements. These are expected to be set after the extension has been created. All RLS policies and security labels on extension objects will be included in dumps created by pg_dump.

The extension mechanism also has provisions for packaging modification scripts that adjust the definitions of the SQL objects contained in an extension. For example, if version 1.1 of an extension adds one function and changes the body of another function compared to 1.0, the extension author can provide an update script that makes just those two changes. The ALTER EXTENSION UPDATE command can then be used to apply these changes and track which version of the extension is actually installed in a given database.

The kinds of SQL objects that can be members of an extension are shown in the description of ALTER EXTENSION. Notably, objects that are database-cluster-wide, such as databases, roles, and tablespaces, cannot be extension members since an extension is only known within one database. (Although an extension script is not prohibited from creating such objects, if it does so they will not be tracked as part of the extension.) Also notice that while a table can be a member of an extension, its subsidiary objects such as indexes are not directly considered members of the extension. Another important point is that schemas can belong to extensions, but not vice versa: an extension as such has an unqualified name and does not exist “within” any schema. The extension’s member objects, however, will belong to schemas whenever appropriate for their object types. It may or may not be appropriate for an extension to own the schema(s) its member objects are within.
36.15.1. Defining Extension Objects

Widely-distributed extensions should assume little about the database they occupy. In particular, unless you issued \texttt{SET search\_path = pg\_temp}, assume each unqualified name could resolve to an object that a malicious user has defined. Beware of constructs that depend on \texttt{search\_path} implicitly: \texttt{IN} and \texttt{CASE expression WHEN} always select an operator using the search path. In their place, \texttt{use OPERATOR(schema.\textasciitilde) ANY} and \texttt{CASE WHEN expression}.

36.15.2. Extension Files

The \texttt{CREATE EXTENSION} command relies on a control file for each extension, which must be named the same as the extension with a suffix of \texttt{.control}, and must be placed in the installation’s \texttt{SHAREDIR/extension} directory. There must also be at least one SQL script file, which follows the naming pattern \texttt{extension--version.sql} (for example, \texttt{foo--1.0.sql} for version 1.0 of extension \textit{foo}). By default, the script file(s) are also placed in the \texttt{SHAREDIR/extension} directory; but the control file can specify a different directory for the script file(s).

The file format for an extension control file is the same as for the \texttt{postgresql.conf} file, namely a list of \texttt{parameter\_name = value} assignments, one per line. Blank lines and comments introduced by \# are allowed. Be sure to quote any value that is not a single word or number.

A control file can set the following parameters:

\begin{verbatim}
directory (string)
   The directory containing the extension’s SQL script file(s). Unless an absolute path is given, the name is relative to the installation’s \texttt{SHAREDIR directory}. The default behavior is equivalent to specifying directory = ‘extension’.

default\_version (string)
   The default version of the extension (the one that will be installed if no version is specified in \texttt{CREATE EXTENSION}). Although this can be omitted, that will result in \texttt{CREATE EXTENSION} failing if no \texttt{VERSION} option appears, so you generally don’t want to do that.

comment (string)
   A comment (any string) about the extension. The comment is applied when initially creating an extension, but not during extension updates (since that might override user-added comments). Alternatively, the extension’s comment can be set by writing a \texttt{COMMENT} command in the script file.

encoding (string)
   The character set encoding used by the script file(s). This should be specified if the script files contain any non-ASCII characters. Otherwise the files will be assumed to be in the database encoding.

module\_pathname (string)
   The value of this parameter will be substituted for each occurrence of \texttt{MODULE\_PATHNAME} in the script file(s). If it is not set, no substitution is made. Typically, this is set to \texttt{$libdir/shared\_library\_name} and then \texttt{MODULE\_PATHNAME} is used in \texttt{CREATE FUNCTION} commands for C-language functions, so that the script files do not need to hard-wire the name of the shared library.
\end{verbatim}
Chapter 36. Extending SQL

requires (string)

A list of names of extensions that this extension depends on, for example requires = ‘foo, bar’. Those extensions must be installed before this one can be installed.

superuser (boolean)

If this parameter is true (which is the default), only superusers can create the extension or update it to a new version. If it is set to false, just the privileges required to execute the commands in the installation or update script are required.

relocatable (boolean)

An extension is relocatable if it is possible to move its contained objects into a different schema after initial creation of the extension. The default is false, i.e. the extension is not relocatable. See Section 36.15.3 for more information.

schema (string)

This parameter can only be set for non-relocatable extensions. It forces the extension to be loaded into exactly the named schema and not any other. The schema parameter is consulted only when initially creating an extension, not during extension updates. See Section 36.15.3 for more information.

In addition to the primary control file extension.control, an extension can have secondary control files named in the style extension--version.control. If supplied, these must be located in the script file directory. Secondary control files follow the same format as the primary control file. Any parameters set in a secondary control file override the primary control file when installing or updating to that version of the extension. However, the parameters directory and default_version cannot be set in a secondary control file.

An extension’s SQL script files can contain any SQL commands, except for transaction control commands (BEGIN, COMMIT, etc) and commands that cannot be executed inside a transaction block (such as VACUUM). This is because the script files are implicitly executed within a transaction block.

An extension’s SQL script files can also contain lines beginning with \echo, which will be ignored (treated as comments) by the extension mechanism. This provision is commonly used to throw an error if the script file is fed to psql rather than being loaded via CREATE EXTENSION (see example script in Section 36.15.6). Without that, users might accidentally load the extension’s contents as “loose” objects rather than as an extension, a state of affairs that’s a bit tedious to recover from.

While the script files can contain any characters allowed by the specified encoding, control files should contain only plain ASCII, because there is no way for PostgreSQL to know what encoding a control file is in. In practice this is only an issue if you want to use non-ASCII characters in the extension’s comment. Recommended practice in that case is to not use the control file comment parameter, but instead use COMMENT ON EXTENSION within a script file to set the comment.

36.15.3. Extension Relocatability

Users often wish to load the objects contained in an extension into a different schema than the extension’s author had in mind. There are three supported levels of relocatability:

- A fully relocatable extension can be moved into another schema at any time, even after it’s been loaded into a database. This is done with the ALTER EXTENSION SET SCHEMA command, which automatically renames all the member objects into the new schema. Normally, this is only possible if the extension contains no internal assumptions about what schema any of its objects are in. Also, the extension’s objects must all be in one schema to begin with (ignoring objects that do not
belong to any schema, such as procedural languages). Mark a fully relocatable extension by setting `relocatable = true` in its control file.

- An extension might be relocatable during installation but not afterwards. This is typically the case if the extension’s script file needs to reference the target schema explicitly, for example in setting `search_path` properties for SQL functions. For such an extension, set `relocatable = false` in its control file, and use `@extschema@` to refer to the target schema in the script file. All occurrences of this string will be replaced by the actual target schema’s name before the script is executed. The user can set the target schema using the `SCHEMA` option of `CREATE EXTENSION`.

- If the extension does not support relocation at all, set `relocatable = false` in its control file, and also set `schema` to the name of the intended target schema. This will prevent use of the `SCHEMA` option of `CREATE EXTENSION`, unless it specifies the same schema named in the control file. This choice is typically necessary if the extension contains internal assumptions about schema names that can’t be replaced by uses of `@extschema@`. The `@extschema@` substitution mechanism is available in this case too, although it is of limited use since the schema name is determined by the control file.

In all cases, the script file will be executed with `search_path` initially set to point to the target schema; that is, `CREATE EXTENSION` does the equivalent of this:

```
SET LOCAL search_path TO @extschema@;
```

This allows the objects created by the script file to go into the target schema. The script file can change `search_path` if it wishes, but that is generally undesirable. `search_path` is restored to its previous setting upon completion of `CREATE EXTENSION`.

The target schema is determined by the `schema` parameter in the control file if that is given, otherwise by the `SCHEMA` option of `CREATE EXTENSION` if that is given, otherwise the current default object creation schema (the first one in the caller’s `search_path`). When the control file `schema` parameter is used, the target schema will be created if it doesn’t already exist, but in the other two cases it must already exist.

If any prerequisite extensions are listed in `requires` in the control file, their target schemas are appended to the initial setting of `search_path`. This allows their objects to be visible to the new extension’s script file.

Although a non-relocatable extension can contain objects spread across multiple schemas, it is usually desirable to place all the objects meant for external use into a single schema, which is considered the extension’s target schema. Such an arrangement works conveniently with the default setting of `search_path` during creation of dependent extensions.

### 36.15.4. Extension Configuration Tables

Some extensions include configuration tables, which contain data that might be added or changed by the user after installation of the extension. Ordinarily, if a table is part of an extension, neither the table’s definition nor its content will be dumped by `pg_dump`. But that behavior is undesirable for a configuration table; any data changes made by the user need to be included in dumps, or the extension will behave differently after a dump and reload.

To solve this problem, an extension’s script file can mark a table or a sequence it has created as a configuration relation, which will cause `pg_dump` to include the table’s or the sequence’s contents (not its definition) in dumps. To do that, call the function `pg_extension_config_dump(regclass, text)` after creating the table or the sequence, for example

```
CREATE TABLE my_config (key text, value text);
```
CREATE SEQUENCE my_config_seq;

SELECT pg_catalog.pg_extension_config_dump('my_config', "");
SELECT pg_catalog.pg_extension_config_dump('my_config_seq', ");

Any number of tables or sequences can be marked this way. Sequences associated with serial or bigserial columns can be marked as well.

When the second argument of `pg_extension_config_dump` is an empty string, the entire contents of the table are dumped by pg_dump. This is usually only correct if the table is initially empty as created by the extension script. If there is a mixture of initial data and user-provided data in the table, the second argument of `pg_extension_config_dump` provides a WHERE condition that selects the data to be dumped. For example, you might do

CREATE TABLE my_config (key text, value text, standard_entry boolean);

SELECT pg_catalog.pg_extension_config_dump('my_config', 'WHERE NOT standard_entry');

and then make sure that standard_entry is true only in the rows created by the extension’s script.

For sequences, the second argument of `pg_extension_config_dump` has no effect.

More complicated situations, such as initially-provided rows that might be modified by users, can be handled by creating triggers on the configuration table to ensure that modified rows are marked correctly.

You can alter the filter condition associated with a configuration table by calling `pg_extension_config_dump` again. (This would typically be useful in an extension update script.) The only way to mark a table as no longer a configuration table is to dissociate it from the extension with ALTER EXTENSION ... DROP TABLE.

Note that foreign key relationships between these tables will dictate the order in which the tables are dumped out by pg_dump. Specifically, pg_dump will attempt to dump the referenced-by table before the referencing table. As the foreign key relationships are set up at CREATE EXTENSION time (prior to data being loaded into the tables) circular dependencies are not supported. When circular dependencies exist, the data will still be dumped out but the dump will not be able to be restored directly and user intervention will be required.

Sequences associated with serial or bigserial columns need to be directly marked to dump their state. Marking their parent relation is not enough for this purpose.

### 36.15.5. Extension Updates

One advantage of the extension mechanism is that it provides convenient ways to manage updates to the SQL commands that define an extension’s objects. This is done by associating a version name or number with each released version of the extension’s installation script. In addition, if you want users to be able to update their databases dynamically from one version to the next, you should provide update scripts that make the necessary changes to go from one version to the next. Update scripts have names following the pattern `extension--old_version--target_version.sql` (for example, `foo--1.0--1.1.sql` contains the commands to modify version 1.0 of extension foo into version 1.1).

Given that a suitable update script is available, the command `ALTER EXTENSION UPDATE` will update an installed extension to the specified new version. The update script is run in the same environment that CREATE EXTENSION provides for installation scripts: in particular, search_path is set up in the same way, and any new objects created by the script are automatically added to the extension.
If an extension has secondary control files, the control parameters that are used for an update script are those associated with the script’s target (new) version.

The update mechanism can be used to solve an important special case: converting a “loose” collection of objects into an extension. Before the extension mechanism was added to PostgreSQL (in 9.1), many people wrote extension modules that simply created assorted unpackaged objects. Given an existing database containing such objects, how can we convert the objects into a properly packaged extension? Dropping them and then doing a plain `CREATE EXTENSION` is one way, but it’s not desirable if the objects have dependencies (for example, if there are table columns of a data type created by the extension). The way to fix this situation is to create an empty extension, then use `ALTER EXTENSION ADD` to attach each pre-existing object to the extension, then finally create any new objects that are in the current extension version but were not in the unpackaged release. `CREATE EXTENSION` supports this case with its `FROM old_version` option, which causes it to not run the normal installation script for the target version, but instead the update script named `extension--old_version--target_version.sql`. The choice of the dummy version name to use as `old_version` is up to the extension author, though `unpackaged` is a common convention. If you have multiple prior versions you need to be able to update into extension style, use multiple dummy version names to identify them.

`ALTER EXTENSION` is able to execute sequences of update script files to achieve a requested update. For example, if only `foo--1.0--1.1.sql` and `foo--1.1--2.0.sql` are available, `ALTER EXTENSION` will apply them in sequence if an update to version 2.0 is requested when 1.0 is currently installed.

PostgreSQL doesn’t assume anything about the properties of version names: for example, it does not know whether 1.1 follows 1.0. It just matches up the available version names and follows the path that requires applying the fewest update scripts. (A version name can actually be any string that doesn’t contain `--` or leading or trailing `-`.)

Sometimes it is useful to provide “downgrade” scripts, for example `foo--1.1--1.0.sql` to allow reverting the changes associated with version 1.1. If you do that, be careful of the possibility that a downgrade script might unexpectedly get applied because it yields a shorter path. The risky case is where there is a “fast path” update script that jumps ahead several versions as well as a downgrade script to the fast path’s start point. It might take fewer steps to apply the downgrade and then the fast path than to move ahead one version at a time. If the downgrade script drops any irreplaceable objects, this will yield undesirable results.

To check for unexpected update paths, use this command:

```sql
SELECT * FROM pg_extension_update_paths('extension_name');
```

This shows each pair of distinct known version names for the specified extension, together with the update path sequence that would be taken to get from the source version to the target version, or `NULL` if there is no available update path. The path is shown in textual form with `--` separators. You can use `regexp_split_to_array(path, '--')` if you prefer an array format.

### 36.15.6. Extension Example

Here is a complete example of an SQL-only extension, a two-element composite type that can store any type of value in its slots, which are named “k” and “v”. Non-text values are automatically coerced to text for storage.

The script file `pair--1.0.sql` looks like this:

```sql
-- complain if script is sourced in psql, rather than via CREATE EXTENSION
```
CREATE TYPE pair AS ( k text, v text );

CREATE OR REPLACE FUNCTION pair(text, text)
RETURNS pair LANGUAGE SQL AS 'SELECT ROW($1, $2)::@extschema@.pair;';

CREATE OPERATOR -> (LEFTARG = text, RIGHTARG = text, PROCEDURE = pair);

-- "SET search_path" is easy to get right, but qualified names perform better.
CREATE OR REPLACE FUNCTION lower(pair)
RETURNS pair LANGUAGE SQL
AS 'SELECT ROW(lower($1.k), lower($1.v))::@extschema@.pair;'
SET search_path = pg_temp;

CREATE OR REPLACE FUNCTION pair_concat(pair, pair)
RETURNS pair LANGUAGE SQL
AS 'SELECT ROW($1.k OPERATOR(pg_catalog.||) $2.k,
             $1.v OPERATOR(pg_catalog.||) $2.v)::@extschema@.pair;';

The control file pair.control looks like this:

# pair extension
comment = 'A key/value pair data type'
default_version = '1.0'
relocatable = false

While you hardly need a makefile to install these two files into the correct directory, you could use a Makefile containing this:

EXTENSION = pair
DATA = pair--1.0.sql
PG_CONFIG = pg_config
PGXS := $(shell $(PG_CONFIG) --pgxs)
include $(PGXS)

This makefile relies on PGXS, which is described in Section 36.16. The command make install will install the control and script files into the correct directory as reported by pg_config.

Once the files are installed, use the CREATE EXTENSION command to load the objects into any particular database.

36.16. Extension Building Infrastructure

If you are thinking about distributing your PostgreSQL extension modules, setting up a portable build system for them can be fairly difficult. Therefore the PostgreSQL installation provides a build infrastructure for extensions, called PGXS, so that simple extension modules can be built simply against an already installed server. PGXS is mainly intended for extensions that include C code, although it can be used for pure-SQL extensions too. Note that PGXS is not intended to be a universal build system framework that can be used to build any software interfacing to PostgreSQL; it simply automates
common build rules for simple server extension modules. For more complicated packages, you might need to write your own build system.

To use the PGXS infrastructure for your extension, you must write a simple makefile. In the makefile, you need to set some variables and include the global PGXS makefile. Here is an example that builds an extension module named isbn_issn, consisting of a shared library containing some C code, an extension control file, a SQL script, and a documentation text file:

```makefile
MODULES = isbn_issn
EXTENSION = isbn_issn
DATA = isbn_issn--1.0.sql
DOCS = README.isbn_issn

PG_CONFIG = pg_config
PGXS := $(shell $(PG_CONFIG) --pgxs)
include $(PGXS)
```

The last three lines should always be the same. Earlier in the file, you assign variables or add custom make rules.

Set one of these three variables to specify what is built:

- **MODULES**
  - list of shared-library objects to be built from source files with same stem (do not include library suffixes in this list)
- **MODULE_big**
  - a shared library to build from multiple source files (list object files in OBJS)
- **PROGRAM**
  - an executable program to build (list object files in OBJS)

The following variables can also be set:

- **EXTENSION**
  - extension name(s); for each name you must provide an extension.control file, which will be installed into prefix/share/extension
- **MODULEDIR**
  - subdirectory of prefix/share into which DATA and DOCS files should be installed (if not set, default is extension if EXTENSION is set, or contrib if not)
- **DATA**
  - random files to install into prefix/share/$MODULEDIR
- **DATA_built**
  - random files to install into prefix/share/$MODULEDIR, which need to be built first
- **DATA_TSEARCH**
  - random files to install under prefix/share/tsearch_data
- **DOCS**
  - random files to install under prefix/doc/$MODULEDIR
- **SCRIPTS**
  - script files (not binaries) to install into prefix/bin
SCRIPTS_built

script files (not binaries) to install into prefix/bin, which need to be built first

REGRESS

list of regression test cases (without suffix), see below

REGRESS_OPTS

additional switches to pass to pg_regress

EXTRA_CLEAN

extra files to remove in make clean

PG_CPPFLAGS

will be prepended to CPPFLAGS

PG_CFLAGS

will be appended to CFLAGS

PG_CXXFLAGS

will be appended to CXXFLAGS

PG_LDFLAGS

will be prepended to LDFLAGS

PG_LIBS

will be added to PROGRAM link line

SHLIB_LINK

will be added to MODULE_big link line

PG_CONFIG

path to pg_config program for the PostgreSQL installation to build against (typically just pg_config to use the first one in your PATH)

Put this makefile as Makefile in the directory which holds your extension. Then you can do make to compile, and then make install to install your module. By default, the extension is compiled and installed for the PostgreSQL installation that corresponds to the first pg_config program found in your PATH. You can use a different installation by setting PG_CONFIG to point to its pg_config program, either within the makefile or on the make command line.

You can also run make in a directory outside the source tree of your extension, if you want to keep the build directory separate. This procedure is also called a VPATH build. Here’s how:

mkdir build_dir
cd build_dir
make -f /path/to/extension/source/tree/Makefile
make -f /path/to/extension/source/tree/Makefile install

Alternatively, you can set up a directory for a VPATH build in a similar way to how it is done for the core code. One way to do this is using the core script config/prep_buildtree. Once this has been done you can build by setting the make variable VPATH like this:

make VPATH=/path/to/extension/source/tree
This procedure can work with a greater variety of directory layouts.

The scripts listed in the \texttt{REGRESS} variable are used for regression testing of your module, which can be invoked by \texttt{make installcheck} after doing \texttt{make install}. For this to work you must have a running PostgreSQL server. The script files listed in \texttt{REGRESS} must appear in a subdirectory named \texttt{sql/} in your extension’s directory. These files must have extension \texttt{.sql}, which must not be included in the \texttt{REGRESS} list in the makefile. For each test there should also be a file containing the expected output in a subdirectory named \texttt{expected/}, with the same stem and extension \texttt{.out}. \texttt{make installcheck} executes each test script with \texttt{psql}, and compares the resulting output to the matching expected file. Any differences will be written to the file \texttt{regression.diffs} in \texttt{diff -c} format. Note that trying to run a test that is missing its expected file will be reported as “trouble”, so make sure you have all expected files.

\textbf{Tip:} The easiest way to create the expected files is to create empty files, then do a test run (which will of course report differences). Inspect the actual result files found in the \texttt{results/} directory, then copy them to \texttt{expected/} if they match what you expect from the test.
Chapter 37. Triggers

This chapter provides general information about writing trigger functions. Trigger functions can be written in most of the available procedural languages, including PL/pgSQL (Chapter 41), PL/Tcl (Chapter 42), PL/Perl (Chapter 43), and PL/Python (Chapter 44). After reading this chapter, you should consult the chapter for your favorite procedural language to find out the language-specific details of writing a trigger in it.

It is also possible to write a trigger function in C, although most people find it easier to use one of the procedural languages. It is not currently possible to write a trigger function in the plain SQL function language.

37.1. Overview of Trigger Behavior

A trigger is a specification that the database should automatically execute a particular function whenever a certain type of operation is performed. Triggers can be attached to tables, views, and foreign tables.

On tables and foreign tables, triggers can be defined to execute either before or after any INSERT, UPDATE, or DELETE operation, either once per modified row, or once per SQL statement. If an INSERT contains an ON CONFLICT DO UPDATE clause, it is possible that the effects of a BEFORE insert trigger and a BEFORE update trigger can both be applied together, if a reference to an EXCLUDED column appears. UPDATE triggers can moreover be set to fire only if certain columns are mentioned in the SET clause of the UPDATE statement. Triggers can also fire for TRUNCATE statements. If a trigger event occurs, the trigger’s function is called at the appropriate time to handle the event. Foreign tables do not support the TRUNCATE statement at all.

On views, triggers can be defined to execute instead of INSERT, UPDATE, or DELETE operations. Such INSTEAD OF triggers are fired once for each row that needs to be modified in the view. It is the responsibility of the trigger’s function to perform the necessary modifications to the view’s underlying base table(s) and, where appropriate, return the modified row as it will appear in the view. Triggers on views can also be defined to execute once per SQL statement, before or after INSERT, UPDATE, or DELETE operations. However, such triggers are fired only if there is also an INSTEAD OF trigger on the view. Otherwise, any statement targeting the view must be rewritten into a statement affecting its underlying base table(s), and then the triggers that will be fired are the ones attached to the base table(s).

The trigger function must be defined before the trigger itself can be created. The trigger function must be declared as a function taking no arguments and returning type trigger. (The trigger function receives its input through a specially-passed TriggerData structure, not in the form of ordinary function arguments.)

Once a suitable trigger function has been created, the trigger is established with CREATE TRIGGER. The same trigger function can be used for multiple triggers.

PostgreSQL offers both per-row triggers and per-statement triggers. With a per-row trigger, the trigger function is invoked once for each row that is affected by the statement that fired the trigger. In contrast, a per-statement trigger is invoked only once when an appropriate statement is executed, regardless of the number of rows affected by that statement. In particular, a statement that affects zero rows will still result in the execution of any applicable per-statement triggers. These two types of triggers are sometimes called row-level triggers and statement-level triggers, respectively. Triggers on TRUNCATE may only be defined at statement level. On views, triggers that fire before or after may only be defined...
at statement level, while triggers that fire instead of an INSERT, UPDATE, or DELETE may only be
defined at row level.

Triggers are also classified according to whether they fire before, after, or instead of the operation. These are referred to as BEFORE triggers, AFTER triggers, and INSTEAD OF triggers respectively. Statement-level BEFORE triggers naturally fire before the statement starts to do anything, while statement-level AFTER triggers fire at the very end of the statement. These types of triggers may be defined on tables or views. Row-level BEFORE triggers fire immediately before a particular row is operated on, while row-level AFTER triggers fire at the end of the statement (but before any statement-level AFTER triggers). These types of triggers may only be defined on tables and foreign tables. Row-level INSTEAD OF triggers may only be defined on views, and fire immediately as each row in the view is identified as needing to be operated on.

If an INSERT contains an ON CONFLICT DO UPDATE clause, it is possible that the effects of all row-level BEFORE INSERT triggers and all row-level BEFORE UPDATE triggers can both be applied in a way that is apparent from the final state of the updated row, if an EXCLUDED column is referenced. There need not be an EXCLUDED column reference for both sets of row-level BEFORE triggers to execute, though. The possibility of surprising outcomes should be considered when there are both BEFORE INSERT and BEFORE UPDATE row-level triggers that both affect a row being inserted/updated (this can still be problematic if the modifications are more or less equivalent if they’re not also idempotent). Note that statement-level UPDATE triggers are executed when ON CONFLICT DO UPDATE is specified, regardless of whether or not any rows were affected by the UPDATE (and regardless of whether the alternative UPDATE path was ever taken). An INSERT with an ON CONFLICT DO UPDATE clause will execute statement-level BEFORE INSERT triggers first, then statement-level BEFORE UPDATE triggers, followed by statement-level AFTER UPDATE triggers and finally statement-level AFTER INSERT triggers.

Trigger functions invoked by per-statement triggers should always return NULL. Trigger functions invoked by per-row triggers can return a table row (a value of type HeapTuple) to the calling executor, if they choose. A row-level trigger fired before an operation has the following choices:

- It can return NULL to skip the operation for the current row. This instructs the executor to not perform the row-level operation that invoked the trigger (the insertion, modification, or deletion of a particular table row).
- For row-level INSERT and UPDATE triggers only, the returned row becomes the row that will be inserted or will replace the row being updated. This allows the trigger function to modify the row being inserted or updated.

A row-level BEFORE trigger that does not intend to cause either of these behaviors must be careful to return as its result the same row that was passed in (that is, the NEW row for INSERT and UPDATE triggers, the OLD row for DELETE triggers).

A row-level INSTEAD OF trigger should either return NULL to indicate that it did not modify any data from the view’s underlying base tables, or it should return the view row that was passed in (the NEW row for INSERT and UPDATE operations, or the OLD row for DELETE operations). A nonnull return value is used to signal that the trigger performed the necessary data modifications in the view. This will cause the count of the number of rows affected by the command to be incremented. For INSERT and UPDATE operations only, the trigger may modify the NEW row before returning it. This will change the data returned by INSERT RETURNING or UPDATE RETURNING, and is useful when the view will not show exactly the same data that was provided.

The return value is ignored for row-level triggers fired after an operation, and so they can return NULL.

If more than one trigger is defined for the same event on the same relation, the triggers will be
fired in alphabetical order by trigger name. In the case of BEFORE and INSTEAD OF triggers, the
possibly-modified row returned by each trigger becomes the input to the next trigger. If any BEFORE or INSTEAD OF trigger returns NULL, the operation is abandoned for that row and subsequent triggers are not fired (for that row).

A trigger definition can also specify a Boolean WHEN condition, which will be tested to see whether the trigger should be fired. In row-level triggers the WHEN condition can examine the old and/or new values of columns of the row. (Statement-level triggers can also have WHEN conditions, although the feature is not so useful for them.) In a BEFORE trigger, the WHEN condition is evaluated just before the function is or would be executed, so using WHEN is not materially different from testing the same condition at the beginning of the trigger function. However, in an AFTER trigger, the WHEN condition is evaluated just after the row update occurs, and it determines whether an event is queued to fire the trigger at the end of statement. So when an AFTER trigger’s WHEN condition does not return true, it is not necessary to queue an event nor to re-fetch the row at end of statement. This can result in significant speedups in statements that modify many rows, if the trigger only needs to be fired for a few of the rows. INSTEAD OF triggers do not support WHEN conditions.

Typically, row-level BEFORE triggers are used for checking or modifying the data that will be inserted or updated. For example, a BEFORE trigger might be used to insert the current time into a timestamp column, or to check that two elements of the row are consistent. Row-level AFTER triggers are most sensibly used to propagate the updates to other tables, or make consistency checks against other tables. The reason for this division of labor is that an AFTER trigger can be certain it is seeing the final value of the row, while a BEFORE trigger cannot; there might be other BEFORE triggers firing after it. If you have no specific reason to make a trigger BEFORE or AFTER, the BEFORE case is more efficient, since the information about the operation doesn’t have to be saved until end of statement.

If a trigger function executes SQL commands then these commands might fire triggers again. This is known as cascading triggers. There is no direct limitation on the number of cascade levels. It is possible for cascades to cause a recursive invocation of the same trigger; for example, an INSERT trigger might execute a command that inserts an additional row into the same table, causing the INSERT trigger to be fired again. It is the trigger programmer’s responsibility to avoid infinite recursion in such scenarios.

When a trigger is being defined, arguments can be specified for it. The purpose of including arguments in the trigger definition is to allow different triggers with similar requirements to call the same function. As an example, there could be a generalized trigger function that takes as its arguments two column names and puts the current user in one and the current time stamp in the other. Properly written, this trigger function would be independent of the specific table it is triggering on. So the same function could be used for INSERT events on any table with suitable columns, to automatically track creation of records in a transaction table for example. It could also be used to track last-update events if defined as an UPDATE trigger.

Each programming language that supports triggers has its own method for making the trigger input data available to the trigger function. This input data includes the type of trigger event (e.g., INSERT or UPDATE) as well as any arguments that were listed in CREATE TRIGGER. For a row-level trigger, the input data also includes the NEW row for INSERT and UPDATE triggers, and/or the OLD row for UPDATE and DELETE triggers. Statement-level triggers do not currently have any way to examine the individual row(s) modified by the statement.

37.2. Visibility of Data Changes

If you execute SQL commands in your trigger function, and these commands access the table that the trigger is for, then you need to be aware of the data visibility rules, because they determine whether
these SQL commands will see the data change that the trigger is fired for. Briefly:

- Statement-level triggers follow simple visibility rules: none of the changes made by a statement are visible to statement-level **BEFORE** triggers, whereas all modifications are visible to statement-level **AFTER** triggers.
- The data change (insertion, update, or deletion) causing the trigger to fire is naturally not visible to SQL commands executed in a row-level **BEFORE** trigger, because it hasn’t happened yet.
- However, SQL commands executed in a row-level **BEFORE** trigger will see the effects of data changes for rows previously processed in the same outer command. This requires caution, since the ordering of these change events is not in general predictable; a SQL command that affects multiple rows can visit the rows in any order.
- Similarly, a row-level **INSTEAD OF** trigger will see the effects of data changes made by previous firings of **INSTEAD OF** triggers in the same outer command.
- When a row-level **AFTER** trigger is fired, all data changes made by the outer command are already complete, and are visible to the invoked trigger function.

If your trigger function is written in any of the standard procedural languages, then the above statements apply only if the function is declared **VOLATILE**. Functions that are declared **STABLE** or **IMMUTABLE** will not see changes made by the calling command in any case.

Further information about data visibility rules can be found in Section 45.4. The example in Section 37.4 contains a demonstration of these rules.

### 37.3. Writing Trigger Functions in C

This section describes the low-level details of the interface to a trigger function. This information is only needed when writing trigger functions in C. If you are using a higher-level language then these details are handled for you. In most cases you should consider using a procedural language before writing your triggers in C. The documentation of each procedural language explains how to write a trigger in that language.

Trigger functions must use the “version 1” function manager interface.

When a function is called by the trigger manager, it is not passed any normal arguments, but it is passed a “context” pointer pointing to a **TriggerData** structure. C functions can check whether they were called from the trigger manager or not by executing the macro:

```
CALLED_AS_TRIGGER(fcinfo)
```

which expands to:

```
((fcinfo)->context != NULL && IsA((fcinfo)->context, TriggerData))
```

If this returns true, then it is safe to cast `fcinfo->context` to type `TriggerData *` and make use of the pointed-to `TriggerData` structure. The function must not alter the `TriggerData` structure or any of the data it points to.

**struct TriggerData is defined in commands/trigger.h:**

```c
typedef struct TriggerData
```


Chapter 37. Triggers

NodeTag type;
TriggerEvent tg_event;
Relation tg_relation;
HeapTuple tg_trigtuple;
HeapTuple tg_newtuple;
Trigger *tg_trigger;
Buffer tg_trigtuplebuf;
Buffer tg_newtuplebuf;
} TriggerData;

where the members are defined as follows:

type

Always T_TriggerData.
tg_event

Describes the event for which the function is called. You can use the following macros to examine tg_event:

TRIGGER_FIRED_BEFORE(tg_event)

Returns true if the trigger fired before the operation.

TRIGGER_FIRED_AFTER(tg_event)

Returns true if the trigger fired after the operation.

TRIGGER_FIRED_INSTEAD(tg_event)

Returns true if the trigger fired instead of the operation.

TRIGGER_FIRED_FOR_ROW(tg_event)

Returns true if the trigger fired for a row-level event.

TRIGGER_FIRED_FOR_STATEMENT(tg_event)

Returns true if the trigger fired for a statement-level event.

TRIGGER_FIRED_BY_INSERT(tg_event)

Returns true if the trigger was fired by an INSERT command.

TRIGGER_FIRED_BY_UPDATE(tg_event)

Returns true if the trigger was fired by an UPDATE command.

TRIGGER_FIRED_BY_DELETE(tg_event)

Returns true if the trigger was fired by a DELETE command.

TRIGGER_FIRED_BY_TRUNCATE(tg_event)

Returns true if the trigger was fired by a TRUNCATE command.

tg_relation

A pointer to a structure describing the relation that the trigger fired for. Look at utils/rel.h for details about this structure. The most interesting things are tg_relation->rd_att (descriptor of the relation tuples) and tg_relation->rd_rel->relname (relation name; the type is not char* but NameData; use SPI_getrelname(tg_relation) to get a char* if you need a copy of the name).
Chapter 37. Triggers

tg_trigtuple

A pointer to the row for which the trigger was fired. This is the row being inserted, updated, or deleted. If this trigger was fired for an INSERT or DELETE then this is what you should return from the function if you don’t want to replace the row with a different one (in the case of INSERT) or skip the operation. For triggers on foreign tables, values of system columns herein are unspecified.

tg_newtuple

A pointer to the new version of the row, if the trigger was fired for an UPDATE, and NULL if it is for an INSERT or a DELETE. This is what you have to return from the function if the event is an UPDATE and you don’t want to replace this row by a different one or skip the operation. For triggers on foreign tables, values of system columns herein are unspecified.

tg_trigger

A pointer to a structure of type Trigger, defined in `utils/reltrigger.h`:

```c
typedef struct Trigger {
    Oid         tgoid;
    char        *tgname;
    Oid         tgfoid;
    int16       tgtype;
    char        tgenabled;
    bool        tgisinternal;
    Oid         tgconstrrelid;
    Oid         tgconstrindid;
    Oid         tgconstraint;
    bool        tgdeferrable;
    bool        tginitdeferred;
    int16       tgnargs;
    int16       tgnattr;
    int16       *tgattr;
    char        **tgargs;
    char        *tgqual;
} Trigger;
```

where tgname is the trigger’s name, tgnargs is the number of arguments in tgargs, and tgargs is an array of pointers to the arguments specified in the CREATE TRIGGER statement. The other members are for internal use only.

tg_trigtuplebuf

The buffer containing tg_trigtuple, or InvalidBuffer if there is no such tuple or it is not stored in a disk buffer.

tg_newtuplebuf

The buffer containing tg_newtuple, or InvalidBuffer if there is no such tuple or it is not stored in a disk buffer.

A trigger function must return either a HeapTuple pointer or a NULL pointer (not an SQL null value, that is, do not set isNull true). Be careful to return either tg_trigtuple or tg_newtuple, as appropriate, if you don’t want to modify the row being operated on.
37.4. A Complete Trigger Example

Here is a very simple example of a trigger function written in C. (Examples of triggers written in procedural languages can be found in the documentation of the procedural languages.)

The function `trigf` reports the number of rows in the table `ttest` and skips the actual operation if the command attempts to insert a null value into the column `x`. (So the trigger acts as a not-null constraint but doesn’t abort the transaction.)

First, the table definition:

```
CREATE TABLE ttest (  
    x integer  
);
```

This is the source code of the trigger function:

```
#include "postgres.h"
#include "executor/spi.h" /* this is what you need to work with SPI */
#include "commands/trigger.h" /* ... triggers ... */
#include "utils/rel.h" /* ... and relations */

PG_MODULE_MAGIC;

PG_FUNCTION_INFO_V1(trigf);

Datum trigf(PG_FUNCTION_ARGS) {
    TriggerData *trigdata = (TriggerData *) fcinfo->context;
    TupleDesc tupdesc;
    HeapTuple rettuple;
    char *when;
    bool checknull = false;
    bool isnull;
    int ret, i;

    /* make sure it’s called as a trigger at all */
    if (!CALLED_AS_TRIGGER(fcinfo))
       elog(ERROR, "trigf: not called by trigger manager");

    /* tuple to return to executor */
    if (TRIGGER_FIRED_BY_UPDATE(trigdata->tg_event))
        rettuple = trigdata->tg_newtuple;
    else
        rettuple = trigdata->tg_trigtuple;

    /* check for null values */
    if (!TRIGGER_FIRED_BY_DELETE(trigdata->tg_event)  
        && TRIGGER_FIRED_BEFORE(trigdata->tg_event))
        checknull = true;

    if (TRIGGER_FIRED_BEFORE(trigdata->tg_event))
        when = "before";
    else
        when = "after ";

    /* ... */
}
```
tupdesc = trigdata->tg_relation->rd_att;

/* connect to SPI manager */
if ((ret = SPI_connect()) < 0)
    elog(ERROR, "trigf (fired %s): SPI_connect returned %d", when, ret);

/* get number of rows in table */
ret = SPI_exec("SELECT count(*) FROM ttest", 0);

if (ret < 0)
    elog(ERROR, "trigf (fired %s): SPI_exec returned %d", when, ret);

/* count(*) returns int8, so be careful to convert */
i = DatumGetInt64(SPI_getbinval(SPI_tuptable->vals[0],
    SPI_tuptable->tupdesc,
    1,
    &isnull));

eilog (INFO, "trigf (fired %s): there are %d rows in ttest", when, i);

SPI_finish();
if (checknull)
{
    SPI_getbinval(rettuple, tupdesc, 1, &isnull);
    if (isnull)
        rettuple = NULL;
}

return PointerGetDatum(rettuple);

After you have compiled the source code (see Section 36.9.6), declare the function and the triggers:

CREATE FUNCTION trigf() RETURNS trigger
    AS 'filename'
    LANGUAGE C;

CREATE TRIGGER tbefore BEFORE INSERT OR UPDATE OR DELETE ON ttest
    FOR EACH ROW EXECUTE PROCEDURE trigf();

CREATE TRIGGER tafter AFTER INSERT OR UPDATE OR DELETE ON ttest
    FOR EACH ROW EXECUTE PROCEDURE trigf();

Now you can test the operation of the trigger:

=> INSERT INTO ttest VALUES (NULL);
INFO: trigf (fired before): there are 0 rows in ttest
INSERT 0 0
-- Insertion skipped and AFTER trigger is not fired

=> SELECT * FROM ttest;
x
---
(0 rows)

=> INSERT INTO ttest VALUES (1);
INFO: trigf (fired before): there are 0 rows in ttest
INFO: trigf (fired after): there are 1 rows in ttest

remember what we said about visibility.

INSERT 167793 1
vac=> SELECT * FROM ttest;
x
---
1
(1 row)

=> INSERT INTO ttest SELECT x * 2 FROM ttest;
INFO: trigf (fired before): there are 1 rows in ttest
INFO: trigf (fired after): there are 2 rows in ttest

remember what we said about visibility.

INSERT 167794 1
=> SELECT * FROM ttest;
x
---
1
2
(2 rows)

=> UPDATE ttest SET x = NULL WHERE x = 2;
INFO: trigf (fired before): there are 2 rows in ttest
UPDATE 0
=> UPDATE ttest SET x = 4 WHERE x = 2;
INFO: trigf (fired before): there are 2 rows in ttest
INFO: trigf (fired after): there are 2 rows in ttest
UPDATE 1
vac=> SELECT * FROM ttest;
x
---
1
4
(2 rows)

=> DELETE FROM ttest;
INFO: trigf (fired before): there are 2 rows in ttest
INFO: trigf (fired before): there are 1 rows in ttest
INFO: trigf (fired after): there are 0 rows in ttest
INFO: trigf (fired after): there are 0 rows in ttest

remember what we said about visibility.

DELETE 2
=> SELECT * FROM ttest;
x
---
(0 rows)

There are more complex examples in src/test/regress/regress.c and in spi.
Chapter 38. Event Triggers

To supplement the trigger mechanism discussed in Chapter 37, PostgreSQL also provides event triggers. Unlike regular triggers, which are attached to a single table and capture only DML events, event triggers are global to a particular database and are capable of capturing DDL events.

Like regular triggers, event triggers can be written in any procedural language that includes event trigger support, or in C, but not in plain SQL.

38.1. Overview of Event Trigger Behavior

An event trigger fires whenever the event with which it is associated occurs in the database in which it is defined. Currently, the only supported events are `ddl_command_start`, `ddl_command_end`, `table_rewrite` and `sql_drop`. Support for additional events may be added in future releases.

The `ddl_command_start` event occurs just before the execution of a `CREATE`, `ALTER`, `DROP`, `SECURITY LABEL`, `COMMENT`, `GRANT` or `REVOKE` command. No check whether the affected object exists or doesn’t exist is performed before the event trigger fires. As an exception, however, this event does not occur for DDL commands targeting shared objects — databases, roles, and tablespaces — or for commands targeting event triggers themselves. The event trigger mechanism does not support these object types. `ddl_command_start` also occurs just before the execution of a `SELECT INTO` command, since this is equivalent to `CREATE TABLE AS`.

The `ddl_command_end` event occurs just after the execution of this same set of commands. To obtain more details on the DDL operations that took place, use the set-returning function `pg_event_trigger_ddl_commands()` from the `ddl_command_end` event trigger code (see Section 9.28). Note that the trigger fires after the actions have taken place (but before the transaction commits), and thus the system catalogs can be read as already changed.

The `sql_drop` event occurs just before the `ddl_command_end` event trigger for any operation that drops database objects. To list the objects that have been dropped, use the set-returning function `pg_event_trigger_dropped_objects()` from the `sql_drop` event trigger code (see Section 9.28). Note that the trigger is executed after the objects have been deleted from the system catalogs, so it’s not possible to look them up anymore.

The `table_rewrite` event occurs just before a table is rewritten by some actions of the commands `ALTER TABLE` and `ALTER TYPE`. While other control statements are available to rewrite a table, like `CLUSTER` and `VACUUM`, the `table_rewrite` event is not triggered by them.

Event triggers (like other functions) cannot be executed in an aborted transaction. Thus, if a DDL command fails with an error, any associated `ddl_command_end` triggers will not be executed. Conversely, if a `ddl_command_start` trigger fails with an error, no further event triggers will fire, and no attempt will be made to execute the command itself. Similarly, if a `ddl_command_end` trigger fails with an error, the effects of the DDL statement will be rolled back, just as they would be in any other case where the containing transaction aborts.

For a complete list of commands supported by the event trigger mechanism, see Section 38.2.

Event triggers are created using the command `CREATE EVENT TRIGGER`. In order to create an event trigger, you must first create a function with the special return type `event_trigger`. This function need not (and may not) return a value; the return type serves merely as a signal that the function is to be invoked as an event trigger.

If more than one event trigger is defined for a particular event, they will fire in alphabetical order by trigger name.
A trigger definition can also specify a \textit{WHEN} condition so that, for example, a \texttt{ddl\_command\_start} trigger can be fired only for particular commands which the user wishes to intercept. A common use of such triggers is to restrict the range of DDL operations which users may perform.

### 38.2. Event Trigger Firing Matrix

Table 38-1 lists all commands for which event triggers are supported.

#### Table 38-1. Event Trigger Support by Command Tag

<table>
<thead>
<tr>
<th>Command Tag</th>
<th>\texttt{ddl_command}</th>
<th>\texttt{alter_command}</th>
<th>\texttt{sql_drop}</th>
<th>\texttt{table_rewrite}</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>ALTER AGGREGATE</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>ALTER COLLATION</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
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<tr>
<td>ALTER CONVERSION</td>
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<td>X</td>
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<td>-</td>
<td></td>
</tr>
<tr>
<td>ALTER DOMAIN</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>ALTER DEFAULT PRIVILEGES</td>
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<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>ALTER EXTENSION</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>ALTER FOREIGN DATA WRAPPER</td>
<td>X</td>
<td>X</td>
<td>-</td>
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<td></td>
</tr>
<tr>
<td>ALTER FOREIGN TABLE</td>
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<tr>
<td>ALTER FUNCTION</td>
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<tr>
<td>ALTER LANGUAGE</td>
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<tr>
<td>ALTER LARGE OBJECT</td>
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</tr>
<tr>
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<tr>
<td>Command Tag</td>
<td>ddl_command</td>
<td>sbtr_command</td>
<td>smpl_drop</td>
<td>table_rewrite</td>
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<td>ALTER TEXT SEARCH DICTIONARY</td>
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<td>ALTER TYPE</td>
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<tr>
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Chapter 38. Event Triggers
## Chapter 38. Event Triggers

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## Chapter 38. Event Triggers

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<td>DROP FUNCTION</td>
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<td>DROP INDEX</td>
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## Chapter 38. Event Triggers

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<thead>
<tr>
<th>Command Tag</th>
<th>ddl_command</th>
<th>short_command</th>
<th>script_drop</th>
<th>table_rewrite</th>
<th>Notes</th>
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<td>DROP OPERATOR CLASS</td>
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<td>DROP OPERATOR FAMILY</td>
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<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP SCHEMA</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP SEQUENCE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP SERVER</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TABLE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TEXT SEARCH CONFIGURATION</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TEXT SEARCH DICTIONARY</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TEXT SEARCH PARSER</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TEXT SEARCH TEMPLATE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TRIGGER</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP TYPE</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP USER MAPPING</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>DROP VIEW</td>
<td>X</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>GRANT</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td>Only for local objects</td>
</tr>
</tbody>
</table>
### 38.3. Writing Event Trigger Functions in C

This section describes the low-level details of the interface to an event trigger function. This information is only needed when writing event trigger functions in C. If you are using a higher-level language then these details are handled for you. In most cases you should consider using a procedural language before writing your event triggers in C. The documentation of each procedural language explains how to write an event trigger in that language.

Event trigger functions must use the “version 1” function manager interface.

When a function is called by the event trigger manager, it is not passed any normal arguments, but it is passed a “context” pointer pointing to a `EventTriggerData` structure. C functions can check whether they were called from the event trigger manager or not by executing the macro:

```c
CALLED_AS_EVENT_TRIGGER(fcinfo)
```

which expands to:

```c
((fcinfo)->context != NULL && IsA((fcinfo)->context, EventTriggerData))
```

If this returns true, then it is safe to cast `fcinfo->context` to type `EventTriggerData *` and make use of the pointed-to `EventTriggerData` structure. The function must not alter the `EventTriggerData` structure or any of the data it points to.

The `EventTriggerData` structure is defined in `commands/event_trigger.h`:

```c
typedef struct EventTriggerData {
    NodeTag type;
    const char *event; /* event name */
    Node *parsetree; /* parse tree */
    const char *tag; /* command tag */
} EventTriggerData;
```

where the members are defined as follows:

<table>
<thead>
<tr>
<th>Command Tag</th>
<th>ddl_command</th>
<th>ddl_tag</th>
<th>sql_drop</th>
<th>table_rewrite</th>
<th>Notes</th>
</tr>
</thead>
<tbody>
<tr>
<td>IMPORT</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
<tr>
<td>FOREIGN</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SCHEMA</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>REFRESH</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td>Only for local objects</td>
</tr>
<tr>
<td>MATERIALIZED VIEW</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>REVOKE</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td>Only for local objects</td>
</tr>
<tr>
<td>SECURITY</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>LABEL</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>SELECT INTO</td>
<td>X</td>
<td>X</td>
<td>-</td>
<td>-</td>
<td></td>
</tr>
</tbody>
</table>

---

**Chapter 38. Event Triggers**
Chapter 38. Event Triggers

type

Always T_EventTriggerData.

event

Describes the event for which the function is called, one of "ddl_command_start", "ddl_command_end", "sql_drop", "table_rewrite". See Section 38.1 for the meaning of these events.

parsetree

A pointer to the parse tree of the command. Check the PostgreSQL source code for details. The parse tree structure is subject to change without notice.

tag

The command tag associated with the event for which the event trigger is run, for example "CREATE FUNCTION".

An event trigger function must return a NULL pointer (not an SQL null value, that is, do not set isNull true).

38.4. A Complete Event Trigger Example

Here is a very simple example of an event trigger function written in C. (Examples of triggers written in procedural languages can be found in the documentation of the procedural languages.)

The function noddl raises an exception each time it is called. The event trigger definition associated the function with the ddl_command_start event. The effect is that all DDL commands (with the exceptions mentioned in Section 38.1) are prevented from running.

This is the source code of the trigger function:

```c
#include "postgres.h"
#include "commands/event_trigger.h"

PG_MODULE_MAGIC;

PG_FUNCTION_INFO_V1(noddl);

Datum noddl(PG_FUNCTION_ARGS)
{
    EventTriggerData *trigdata;
    if (!CALLED_AS_EVENT_TRIGGER(fcinfo)) /* internal error */
        elog(ERROR, "not fired by event trigger manager");
    trigdata = (EventTriggerData *) fcinfo->context;
    ereport(ERROR,
        {errcode(ERRCODE_INSUFFICIENT_PRIVILEGE),
        errmsg("command \"%s\" denied", trigdata->tag)});
    PG_RETURN_NULL();
```
After you have compiled the source code (see Section 36.9.6), declare the function and the triggers:

```sql
CREATE FUNCTION noddl() RETURNS event_trigger
  AS 'noddl' LANGUAGE C;
CREATE EVENT TRIGGER noddl ON ddl_command_start
  EXECUTE PROCEDURE noddl();
```

Now you can test the operation of the trigger:

```sql
=# \d
```

<table>
<thead>
<tr>
<th>List of event triggers</th>
</tr>
</thead>
<tbody>
<tr>
<td>Name</td>
</tr>
<tr>
<td>------</td>
</tr>
<tr>
<td>noddl</td>
</tr>
</tbody>
</table>

```sql
=# CREATE TABLE foo(id serial);
ERROR: command "CREATE TABLE" denied
```

In this situation, in order to be able to run some DDL commands when you need to do so, you have to either drop the event trigger or disable it. It can be convenient to disable the trigger for only the duration of a transaction:

```sql
BEGIN;
ALTER EVENT TRIGGER noddl DISABLE;
CREATE TABLE foo (id serial);
ALTER EVENT TRIGGER noddl ENABLE;
COMMIT;
```

(Recall that DDL commands on event triggers themselves are not affected by event triggers.)

### 38.5. A Table Rewrite Event Trigger Example

Thanks to the `table_rewrite` event, it is possible to implement a table rewriting policy only allowing the rewrite in maintenance windows.

Here’s an example implementing such a policy.

```sql
CREATE OR REPLACE FUNCTION no_rewrite()
  RETURNS event_trigger
  LANGUAGE plpgsql AS
$$
---
--- Implement local Table Rewriting policy:
--- public.foo is not allowed rewriting, ever
--- other tables are only allowed rewriting between 1am and 6am
--- unless they have more than 100 blocks
$$
```
DECLARE
  table_oid oid := pg_event_trigger_table_rewrite_oid();
  current_hour integer := extract('hour' from current_time);
  pages integer;
  max_pages integer := 100;
BEGIN
  IF pg_event_trigger_table_rewrite_oid() = 'public.foo'::regclass
  THEN
    RAISE EXCEPTION 'you’re not allowed to rewrite the table %',
    table_oid::regclass;
  END IF;

  SELECT INTO pages relpages FROM pg_class WHERE oid = table_oid;
  IF pages > max_pages
  THEN
    RAISE EXCEPTION 'rewrites only allowed for table with less than % pages',
    max_pages;
  END IF;

  IF current_hour NOT BETWEEN 1 AND 6
  THEN
    RAISE EXCEPTION 'rewrites only allowed between 1am and 6am';
  END IF;
END;
$$;

CREATE EVENT TRIGGER no_rewrite_allowed
  ON table_rewrite
  EXECUTE PROCEDURE no_rewrite();
Chapter 39. The Rule System

This chapter discusses the rule system in PostgreSQL. Production rule systems are conceptually simple, but there are many subtle points involved in actually using them.

Some other database systems define active database rules, which are usually stored procedures and triggers. In PostgreSQL, these can be implemented using functions and triggers as well.

The rule system (more precisely speaking, the query rewrite rule system) is totally different from stored procedures and triggers. It modifies queries to take rules into consideration, and then passes the modified query to the query planner for planning and execution. It is very powerful, and can be used for many things such as query language procedures, views, and versions. The theoretical foundations and the power of this rule system are also discussed in On Rules, Procedures, Caching and Views in Database Systems and A Unified Framework for Version Modeling Using Production Rules in a Database System.

39.1. The Query Tree

To understand how the rule system works it is necessary to know when it is invoked and what its input and results are.

The rule system is located between the parser and the planner. It takes the output of the parser, one query tree, and the user-defined rewrite rules, which are also query trees with some extra information, and creates zero or more query trees as result. So its input and output are always things the parser itself could have produced and thus, anything it sees is basically representable as an SQL statement.

Now what is a query tree? It is an internal representation of an SQL statement where the single parts that it is built from are stored separately. These query trees can be shown in the server log if you set the configuration parameters debug_print_parse, debug_print_rewritten, or debug_print_plan. The rule actions are also stored as query trees, in the system catalog pg_rewrite. They are not formatted like the log output, but they contain exactly the same information.

Reading a raw query tree requires some experience. But since SQL representations of query trees are sufficient to understand the rule system, this chapter will not teach how to read them.

When reading the SQL representations of the query trees in this chapter it is necessary to be able to identify the parts the statement is broken into when it is in the query tree structure. The parts of a query tree are

the command type

This is a simple value telling which command (SELECT, INSERT, UPDATE, DELETE) produced the query tree.

the range table

The range table is a list of relations that are used in the query. In a SELECT statement these are the relations given after the FROM key word.

Every range table entry identifies a table or view and tells by which name it is called in the other parts of the query. In the query tree, the range table entries are referenced by number rather than by name, so here it doesn’t matter if there are duplicate names as it would in an SQL statement. This can happen after the range tables of rules have been merged in. The examples in this chapter will not have this situation.
the result relation

This is an index into the range table that identifies the relation where the results of the query go. SELECT queries don’t have a result relation. (The special case of SELECT INTO is mostly identical to CREATE TABLE followed by INSERT ... SELECT, and is not discussed separately here.) For INSERT, UPDATE, and DELETE commands, the result relation is the table (or view!) where the changes are to take effect.

the target list

The target list is a list of expressions that define the result of the query. In the case of a SELECT, these expressions are the ones that build the final output of the query. They correspond to the expressions between the key words SELECT and FROM. (* is just an abbreviation for all the column names of a relation. It is expanded by the parser into the individual columns, so the rule system never sees it.) DELETE commands don’t need a normal target list because they don’t produce any result. Instead, the planner adds a special CTID entry to the empty target list, to allow the executor to find the row to be deleted. (CTID is added when the result relation is an ordinary table. If it is a view, a whole-row variable is added instead, by the rule system, as described in Section 39.2.4.) For INSERT commands, the target list describes the new rows that should go into the result relation. It consists of the expressions in the VALUES clause or the ones from the SELECT clause in INSERT ... SELECT. The first step of the rewrite process adds target list entries for any columns that were not assigned to by the original command but have defaults. Any remaining columns (with neither a given value nor a default) will be filled in by the planner with a constant null expression.

For UPDATE commands, the target list describes the new rows that should replace the old ones. In the rule system, it contains just the expressions from the \texttt{SET column = expression} part of the command. The planner will handle missing columns by inserting expressions that copy the values from the old row into the new one. Just as for DELETE, a CTID or whole-row variable is added so that the executor can identify the old row to be updated.

Every entry in the target list contains an expression that can be a constant value, a variable pointing to a column of one of the relations in the range table, a parameter, or an expression tree made of function calls, constants, variables, operators, etc.

the qualification

The query’s qualification is an expression much like one of those contained in the target list entries. The result value of this expression is a Boolean that tells whether the operation (INSERT, UPDATE, DELETE, or SELECT) for the final result row should be executed or not. It corresponds to the \texttt{WHERE} clause of an SQL statement.

the join tree

The query’s join tree shows the structure of the \texttt{FROM} clause. For a simple query like SELECT ... FROM a, b, c, the join tree is just a list of the FROM items, because we are allowed to join them in any order. But when JOIN expressions, particularly outer joins, are used, we have to join in the order shown by the joins. In that case, the join tree shows the structure of the JOIN expressions. The restrictions associated with particular JOIN clauses (from \texttt{ON} or \texttt{USING} expressions) are stored as qualification expressions attached to those join-tree nodes. It turns out to be convenient to store the top-level \texttt{WHERE} expression as a qualification attached to the top-level join-tree item, too. So really the join tree represents both the \texttt{FROM} and \texttt{WHERE} clauses of a SELECT.
Chapter 39. The Rule System

39.2. Views and the Rule System

Views in PostgreSQL are implemented using the rule system. In fact, there is essentially no difference between:

CREATE VIEW myview AS SELECT * FROM mytab;

compared against the two commands:

CREATE TABLE myview (same column list as mytab);
CREATE RULE "_RETURN" AS ON SELECT TO myview DO INSTEAD
SELECT * FROM mytab;

because this is exactly what the CREATE VIEW command does internally. This has some side effects. One of them is that the information about a view in the PostgreSQL system catalogs is exactly the same as it is for a table. So for the parser, there is absolutely no difference between a table and a view. They are the same thing: relations.

39.2.1. How SELECT Rules Work

Rules ON SELECT are applied to all queries as the last step, even if the command given is an INSERT, UPDATE or DELETE. And they have different semantics from rules on the other command types in that they modify the query tree in place instead of creating a new one. So SELECT rules are described first.

Currently, there can be only one action in an ON SELECT rule, and it must be an unconditional SELECT action that is INSTEAD. This restriction was required to make rules safe enough to open them for ordinary users, and it restricts ON SELECT rules to act like views.

The examples for this chapter are two join views that do some calculations and some more views using them in turn. One of the two first views is customized later by adding rules for INSERT, UPDATE, and DELETE operations so that the final result will be a view that behaves like a real table with some magic functionality. This is not such a simple example to start from and this makes things harder to get into. But it’s better to have one example that covers all the points discussed step by step rather than having many different ones that might mix up in mind.

The real tables we need in the first two rule system descriptions are these:

CREATE TABLE shoe_data (  
    shoename  text,  -- primary key
    sh_avail  integer,  -- available number of pairs
    slcolor   text,  -- preferred shoelace color
    slminlen  real,  -- minimum shoelace length
    slmaxlen  real,  -- maximum shoelace length
    slunit    text  -- length unit
);

CREATE TABLE shoelace_data (  
...
Chapter 39. The Rule System

```
sl_name text, -- primary key
sl_avail integer, -- available number of pairs
sl_color text, -- shoelace color
sl_len real, -- shoelace length
sl_unit text -- length unit
);

CREATE TABLE unit (  
  un_name text, -- primary key
  un_fact real -- factor to transform to cm
);

As you can see, they represent shoe-store data.

The views are created as:

CREATE VIEW shoe AS  
SELECT sh.shoename,
    sh.sh_avail,
    sh.slcolor,
    sh.slminlen,
    sh.slminlen * un.un_fact AS slminlen_cm,
    sh.slmaxlen,
    sh.slmaxlen * un.un_fact AS slmaxlen_cm,
    sh.slunit  
FROM shoe_data sh, unit un  
WHERE sh.slunit = un.un_name;

CREATE VIEW shoelace AS  
SELECT s.sl_name,
    s.sl_avail,
    s.sl_color,
    s.sl_len,
    s.sl_unit,
    s.sl_len * u.un_fact AS sl_len_cm  
FROM shoelace_data s, unit u  
WHERE s.sl_unit = u.un_name;

CREATE VIEW shoe_ready AS  
SELECT rsh.shoename,
    rsh.sh_avail,
    rsl.sl_name,
    rsl.sl_avail,
    least(rsh.sh_avail, rsl.sl_avail) AS total_avail  
FROM shoe rsh, shoelace rsl  
WHERE rsl.sl_color = rsh.slcolor  
AND rsl.sl_len_cm >= rsh.slminlen_cm  
AND rsl.sl_len_cm <= rsh.slmaxlen_cm;

The CREATE VIEW command for the shoelace view (which is the simplest one we have) will create a relation shoelace and an entry in pg_rewrite that tells that there is a rewrite rule that must be applied whenever the relation shoelace is referenced in a query’s range table. The rule has no rule qualification (discussed later, with the non-SELECT rules, since SELECT rules currently cannot have them) and it is INSTEAD. Note that rule qualifications are not the same as query qualifications. The action of our rule has a query qualification. The action of the rule is one query tree that is a copy of the SELECT statement in the view creation command.

1088
Chapter 39. The Rule System

Note: The two extra range table entries for \texttt{NEW} and \texttt{OLD} that you can see in the \texttt{pg_rewrite} entry aren't of interest for \texttt{SELECT} rules.

Now we populate \texttt{unit}, \texttt{shoe_data} and \texttt{shoelace_data} and run a simple query on a view:

\begin{verbatim}
INSERT INTO unit VALUES ('cm', 1.0);
INSERT INTO unit VALUES ('m', 100.0);
INSERT INTO unit VALUES ('inch', 2.54);

INSERT INTO shoe_data VALUES ('sh1', 2, 'black', 70.0, 90.0, 'cm');
INSERT INTO shoe_data VALUES ('sh2', 0, 'black', 30.0, 40.0, 'inch');
INSERT INTO shoe_data VALUES ('sh3', 4, 'brown', 50.0, 65.0, 'cm');
INSERT INTO shoe_data VALUES ('sh4', 3, 'brown', 40.0, 50.0, 'inch');

INSERT INTO shoelace_data VALUES ('sl1', 5, 'black', 80.0, 'cm');
INSERT INTO shoelace_data VALUES ('sl2', 6, 'black', 100.0, 'cm');
INSERT INTO shoelace_data VALUES ('sl3', 0, 'black', 35.0, 'inch');
INSERT INTO shoelace_data VALUES ('sl4', 8, 'black', 40.0, 'inch');
INSERT INTO shoelace_data VALUES ('sl5', 4, 'brown', 1.0, 'm');
INSERT INTO shoelace_data VALUES ('sl6', 0, 'brown', 0.9, 'm');
INSERT INTO shoelace_data VALUES ('sl7', 7, 'brown', 60, 'cm');
INSERT INTO shoelace_data VALUES ('sl8', 1, 'brown', 40, 'inch');

SELECT * FROM shoelace;
\end{verbatim}

\begin{verbatim}
<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl1</td>
<td>5</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl2</td>
<td>6</td>
<td>black</td>
<td>100</td>
<td>cm</td>
<td>100</td>
</tr>
<tr>
<td>sl7</td>
<td>7</td>
<td>brown</td>
<td>60</td>
<td>cm</td>
<td>60</td>
</tr>
<tr>
<td>sl3</td>
<td>0</td>
<td>brown</td>
<td>35</td>
<td>inch</td>
<td>88.9</td>
</tr>
<tr>
<td>sl4</td>
<td>8</td>
<td>black</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl8</td>
<td>1</td>
<td>brown</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl5</td>
<td>4</td>
<td>brown</td>
<td>1</td>
<td>m</td>
<td>100</td>
</tr>
<tr>
<td>sl6</td>
<td>0</td>
<td>brown</td>
<td>0.9</td>
<td>m</td>
<td>90</td>
</tr>
</tbody>
</table>
\end{verbatim}

This is the simplest \texttt{SELECT} you can do on our views, so we take this opportunity to explain the basics of view rules. The \texttt{SELECT * FROM shoelace} was interpreted by the parser and produced the query tree:

\begin{verbatim}
SELECT shoelace.sl_name, shoelace.sl_avail, 
      shoelace.sl_color, shoelace.sl_len, 
      shoelace.sl_unit, shoelace.sl_len_cm
FROM shoelace shoelace;
\end{verbatim}

and this is given to the rule system. The rule system walks through the range table and checks if there are rules for any relation. When processing the range table entry for \texttt{shoelace} (the only one up to now) it finds the \_\texttt{RETURN} rule with the query tree:

\begin{verbatim}
SELECT s.sl_name, s.sl_avail, 
      s.sl_color, s.sl_len, s.sl_unit, 
      s.sl_len * u.un_fact AS sl_len_cm
FROM shoelace old, shoelace new,
\end{verbatim}
Chapter 39. The Rule System

shoelace_data s, unit u
WHERE s.sl_unit = u.un_name;

To expand the view, the rewriter simply creates a subquery range-table entry containing the rule’s action query tree, and substitutes this range table entry for the original one that referenced the view. The resulting rewritten query tree is almost the same as if you had typed:

SELECT shoelace.sl_name, shoelace.sl_avail,
      shoelace.sl_color, shoelace.sl_len,
      shoelace.sl_unit, shoelace.sl_len_cm
FROM (SELECT s.sl_name,
      s.sl_avail,
      s.sl_color,
      s.sl_len,
      s.sl_unit,
      s.sl_len * u.un_fact AS sl_len_cm
FROM shoelace_data s, unit u
WHERE s.sl_unit = u.un_name) shoelace;

There is one difference however: the subquery’s range table has two extra entries shoelace old and shoelace new. These entries don’t participate directly in the query, since they aren’t referenced by the subquery’s join tree or target list. The rewriter uses them to store the access privilege check information that was originally present in the range-table entry that referenced the view. In this way, the executor will still check that the user has proper privileges to access the view, even though there’s no direct use of the view in the rewritten query.

That was the first rule applied. The rule system will continue checking the remaining range-table entries in the top query (in this example there are no more), and it will recursively check the range-table entries in the added subquery to see if any of them reference views. (But it won’t expand old or new — otherwise we’d have infinite recursion!) In this example, there are no rewrite rules for shoelace_data or unit, so rewriting is complete and the above is the final result given to the planner.

Now we want to write a query that finds out for which shoes currently in the store we have the matching shoelaces (color and length) and where the total number of exactly matching pairs is greater or equal to two.

SELECT * FROM shoe_ready WHERE total_avail >= 2;

The output of the parser this time is the query tree:

SELECT shoe_ready.shoename, shoe_ready.sh_avail,
        shoe_ready.sl_name, shoe_ready.sl_avail,
        shoe_ready.total_avail
FROM shoe_ready WHERE shoe_ready.total_avail >= 2;

The first rule applied will be the one for the shoe_ready view and it results in the query tree:
Chapter 39. The Rule System

SELECT shoe_ready.shoename, shoe_ready.sh_avail, shoe_ready.sl_name, shoe_ready.sl_avail, shoe_ready.total_avail
FROM (SELECT rsh.shoename, rsh.sh_avail, rsl.sl_name, rsl.sl_avail, least(rsh.sh_avail, rsl.sl_avail) AS total_avail
FROM shoe rsh, shoelace rsl
WHERE rsl.sl_color = rsh.slcolor
AND rsl.sl_len_cm >= rsh.slminlen_cm
AND rsl.sl_len_cm <= rsh.slmaxlen_cm) shoe_ready
WHERE shoe_ready.total_avail >= 2;

Similarly, the rules for shoe and shoelace are substituted into the range table of the subquery, leading to a three-level final query tree:

SELECT shoe_ready.shoename, shoe_ready.sh_avail, shoe_ready.sl_name, shoe_ready.sl_avail, shoe_ready.total_avail
FROM (SELECT rsh.shoename, rsh.sh_avail, rsl.sl_name, rsl.sl_avail, least(rsh.sh_avail, rsl.sl_avail) AS total_avail
FROM (SELECT sh.shoename, sh.sh_avail, sh.slcolor, sh.slminlen, sh.slminlen * un.un_fact AS slminlen_cm, sh.slmaxlen, sh.slmaxlen * un.un_fact AS slmaxlen_cm, sh.slunit
FROM shoe_data sh, unit un
WHERE sh.slunit = un.un_name) rsh,
(SELECT s.sl_name, s.sl_avail, s.sl_color, s.sl_len, s.sl_unit, s.sl_len * u.un_fact AS sl_len_cm
FROM shoelace_data s, unit u
WHERE s.sl_unit = u.un_name) rsl
WHERE rsl.sl_color = rsh.slcolor
AND rsl.sl_len_cm >= rsh.slminlen_cm
AND rsl.sl_len_cm <= rsh.slmaxlen_cm) shoe_ready
WHERE shoe_ready.total_avail >= 2;

This might look inefficient, but the planner will collapse this into a single-level query tree by “pulling up” the subqueries, and then it will plan the joins just as if we’d written them out manually. So collapsing the query tree is an optimization that the rewrite system doesn’t have to concern itself with.
Chapter 39. The Rule System

39.2.2. View Rules in Non-SELECT Statements

Two details of the query tree aren’t touched in the description of view rules above. These are the command type and the result relation. In fact, the command type is not needed by view rules, but the result relation may affect the way in which the query rewriter works, because special care needs to be taken if the result relation is a view.

There are only a few differences between a query tree for a SELECT and one for any other command. Obviously, they have a different command type and for a command other than a SELECT, the result relation points to the range-table entry where the result should go. Everything else is absolutely the same. So having two tables t1 and t2 with columns a and b, the query trees for the two statements:

SELECT t2.b FROM t1, t2 WHERE t1.a = t2.a;

UPDATE t1 SET b = t2.b FROM t2 WHERE t1.a = t2.a;

are nearly identical. In particular:

• The range tables contain entries for the tables t1 and t2.
• The target lists contain one variable that points to column b of the range table entry for table t2.
• The qualification expressions compare the columns a of both range-table entries for equality.
• The join trees show a simple join between t1 and t2.

The consequence is, that both query trees result in similar execution plans: They are both joins over the two tables. For the UPDATE the missing columns from t1 are added to the target list by the planner and the final query tree will read as:

UPDATE t1 SET a = t1.a, b = t2.b FROM t2 WHERE t1.a = t2.a;

and thus the executor run over the join will produce exactly the same result set as:

SELECT t1.a, t2.b FROM t1, t2 WHERE t1.a = t2.a;

But there is a little problem in UPDATE: the part of the executor plan that does the join does not care what the results from the join are meant for. It just produces a result set of rows. The fact that one is a SELECT command and the other is an UPDATE is handled higher up in the executor, where it knows that this is an UPDATE, and it knows that this result should go into table t1. But which of the rows that are there has to be replaced by the new row?

To resolve this problem, another entry is added to the target list in UPDATE (and also in DELETE) statements: the current tuple ID (CTID). This is a system column containing the file block number and position in the block for the row. Knowing the table, the CTID can be used to retrieve the original row of t1 to be updated. After adding the CTID to the target list, the query actually looks like:

SELECT t1.a, t2.b, t1.ctid FROM t1, t2 WHERE t1.a = t2.a;

Now another detail of PostgreSQL enters the stage. Old table rows aren’t overwritten, and this is why ROLLBACK is fast. In an UPDATE, the new result row is inserted into the table (after stripping the CTID) and in the row header of the old row, which the CTID pointed to, the cmax and xmax entries are set to the current command counter and current transaction ID. Thus the old row is hidden, and after the transaction commits the vacuum cleaner can eventually remove the dead row.
Knowing all that, we can simply apply view rules in absolutely the same way to any command. There is no difference.

### 39.2.3. The Power of Views in PostgreSQL

The above demonstrates how the rule system incorporates view definitions into the original query tree. In the second example, a simple `SELECT` from one view created a final query tree that is a join of 4 tables (unit was used twice with different names).

The benefit of implementing views with the rule system is, that the planner has all the information about which tables have to be scanned plus the relationships between these tables plus the restrictive qualifications from the views plus the qualifications from the original query in one single query tree.

And this is still the situation when the original query is already a join over views. The planner has to decide which is the best path to execute the query, and the more information the planner has, the better this decision can be. And the rule system as implemented in PostgreSQL ensures, that this is all information available about the query up to that point.

### 39.2.4. Updating a View

What happens if a view is named as the target relation for an `INSERT`, `UPDATE`, or `DELETE`? Doing the substitutions described above would give a query tree in which the result relation points at a subquery range-table entry, which will not work. There are several ways in which PostgreSQL can support the appearance of updating a view, however.

If the subquery selects from a single base relation and is simple enough, the rewriter can automatically replace the subquery with the underlying base relation so that the `INSERT`, `UPDATE`, or `DELETE` is applied to the base relation in the appropriate way. Views that are “simple enough” for this are called *automatically updatable*. For detailed information on the kinds of view that can be automatically updated, see `CREATE VIEW`.

Alternatively, the operation may be handled by a user-provided `INSTEAD OF` trigger on the view. Rewriting works slightly differently in this case. For `INSERT`, the rewriter does nothing at all with the view, leaving it as the result relation for the query. For `UPDATE` and `DELETE`, it’s still necessary to expand the view query to produce the “old” rows that the command will attempt to update or delete. So the view is expanded as normal, but another unexpanded range-table entry is added to the query to represent the view in its capacity as the result relation.

The problem that now arises is how to identify the rows to be updated in the view. Recall that when the result relation is a table, a special CTID entry is added to the target list to identify the physical locations of the rows to be updated. This does not work if the result relation is a view, because a view does not have any CTID, since its rows do not have actual physical locations. Instead, for an `UPDATE` or `DELETE` operation, a special `wholerow` entry is added to the target list, which expands to include all columns from the view. The executor uses this value to supply the “old” row to the `INSTEAD OF` trigger. It is up to the trigger to work out what to update based on the old and new row values.

Another possibility is for the user to define `INSTEAD` rules that specify substitute actions for `INSERT`, `UPDATE`, and `DELETE` commands on a view. These rules will rewrite the command, typically into a command that updates one or more tables, rather than views. That is the topic of Section 39.4.

Note that rules are evaluated first, rewriting the original query before it is planned and executed. Therefore, if a view has `INSTEAD OF` triggers as well as rules on `INSERT`, `UPDATE`, or `DELETE`, then the rules will be evaluated first, and depending on the result, the triggers may not be used at all.
Automatic rewriting of an `INSERT`, `UPDATE`, or `DELETE` query on a simple view is always tried last. Therefore, if a view has rules or triggers, they will override the default behavior of automatically updatable views.

If there are no `INSTEAD` rules or `INSTEAD OF` triggers for the view, and the rewriter cannot automatically rewrite the query as an update on the underlying base relation, an error will be thrown because the executor cannot update a view as such.

### 39.3. Materialized Views

Materialized views in PostgreSQL use the rule system like views do, but persist the results in a table-like form. The main differences between:

```sql
CREATE MATERIALIZED VIEW mymatview AS SELECT * FROM mytab;
```

and:

```sql
CREATE TABLE mymatview AS SELECT * FROM mytab;
```

are that the materialized view cannot subsequently be directly updated and that the query used to create the materialized view is stored in exactly the same way that a view’s query is stored, so that fresh data can be generated for the materialized view with:

```sql
REFRESH MATERIALIZED VIEW mymatview;
```

The information about a materialized view in the PostgreSQL system catalogs is exactly the same as it is for a table or view. So for the parser, a materialized view is a relation, just like a table or a view. When a materialized view is referenced in a query, the data is returned directly from the materialized view, like from a table; the rule is only used for populating the materialized view.

While access to the data stored in a materialized view is often much faster than accessing the underlying tables directly or through a view, the data is not always current; yet sometimes current data is not needed. Consider a table which records sales:

```sql
CREATE TABLE invoice (  
  invoice_no integer PRIMARY KEY,  
  seller_no integer, -- ID of salesperson  
  invoice_date date, -- date of sale  
  invoice_amt numeric(13,2) -- amount of sale
);
```

If people want to be able to quickly graph historical sales data, they might want to summarize, and they may not care about the incomplete data for the current date:

```sql
CREATE MATERIALIZED VIEW sales_summary AS
  SELECT  
    seller_no,  
    invoice_date,  
    sum(invoice_amt)::numeric(13,2) as sales_amt
  FROM invoice
  WHERE invoice_date < CURRENT_DATE
  GROUP BY  
    seller_no,  
    invoice_date
ORDER BY
```

1094
Chapter 39. The Rule System

seller_no,  
invoice_date;

CREATE UNIQUE INDEX sales_summary_seller  
ON sales_summary (seller_no, invoice_date);

This materialized view might be useful for displaying a graph in the dashboard created for salespeople.  
A job could be scheduled to update the statistics each night using this SQL statement:

REFRESH MATERIALIZED VIEW sales_summary;

Another use for a materialized view is to allow faster access to data brought across from a remote  
system through a foreign data wrapper. A simple example using file_fdw is below, with timings,  
but since this is using cache on the local system the performance difference compared to access to a  
remote system would usually be greater than shown here. Notice we are also exploiting the ability to  
put an index on the materialized view, whereas file_fdw does not support indexes; this advantage  
might not apply for other sorts of foreign data access.

Setup:

CREATE EXTENSION file_fdw;
CREATE SERVER local_file FOREIGN DATA WRAPPER file_fdw;
CREATE FOREIGN TABLE words (word text NOT NULL)  
SERVER local_file  
OPTIONS (filename '/usr/share/dict/words');
CREATE MATERIALIZED VIEW wrd AS SELECT * FROM words;
CREATE UNIQUE INDEX wrd_word ON wrd (word);
CREATE EXTENSION pg_trgm;
CREATE INDEX wrd_trgm ON wrd USING gist (word gist_trgm_ops);
VACUUM ANALYZE wrd;

Now let’s spell-check a word. Using file_fdw directly:

SELECT count(*) FROM words WHERE word = 'caterpilier';

<table>
<thead>
<tr>
<th>count</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
</tr>
</tbody>
</table>

(1 row)

With EXPLAIN ANALYZE, we see:

Aggregate (cost=21763.99..21764.00 rows=1 width=0) (actual time=188.180..188.181 rows=1)
   -> Foreign Scan on words (cost=0.00..21761.41 rows=1032 width=0) (actual time=188.177..188.177 rows=0)
      Filter: (word = 'caterpilier'::text)
      Rows Removed by Filter: 479829
      Foreign File: /usr/share/dict/words
      Foreign File Size: 4953699
Planning time: 0.118 ms
Execution time: 188.273 ms

If the materialized view is used instead, the query is much faster:

Aggregate (cost=4.44..4.45 rows=1 width=0) (actual time=0.042..0.042 rows=1 loops=1)
   -> Index Only Scan using wrd_word on wrd (cost=0.42..4.44 rows=1 width=0) (actual time=0.040..0.040)
      Index Cond: (word = 'caterpilier'::text)
Either way, the word is spelled wrong, so let's look for what we might have wanted. Again using file_fdw:

```
SELECT word FROM words ORDER BY word <-> 'caterpiler' LIMIT 10;
```

<table>
<thead>
<tr>
<th>word</th>
</tr>
</thead>
<tbody>
<tr>
<td>cater</td>
</tr>
<tr>
<td>caterpillar</td>
</tr>
<tr>
<td>Caterpillar</td>
</tr>
<tr>
<td>caterpillars</td>
</tr>
<tr>
<td>caterpillar’s</td>
</tr>
<tr>
<td>Caterpillar’s</td>
</tr>
<tr>
<td>caterer</td>
</tr>
<tr>
<td>caterer’s</td>
</tr>
<tr>
<td>caterers</td>
</tr>
<tr>
<td>catered</td>
</tr>
</tbody>
</table>

(10 rows)

If you can tolerate periodic update of the remote data to the local database, the performance benefit can be substantial.

### 39.4. Rules on `INSERT`, `UPDATE`, and `DELETE`

Rules that are defined on `INSERT`, `UPDATE`, and `DELETE` are significantly different from the view rules described in the previous section. First, their `CREATE RULE` command allows more:

- They are allowed to have no action.
- They can have multiple actions.
- They can be `INSTEAD` or `ALSO` (the default).
• The pseudorelations NEW and OLD become useful.
• They can have rule qualifications.

Second, they don’t modify the query tree in place. Instead they create zero or more new query trees and can throw away the original one.

Caution

In many cases, tasks that could be performed by rules on INSERT/UPDATE/DELETE are better done with triggers. Triggers are notationally a bit more complicated, but their semantics are much simpler to understand. Rules tend to have surprising results when the original query contains volatile functions: volatile functions may get executed more times than expected in the process of carrying out the rules. Also, there are some cases that are not supported by these types of rules at all, notably including WITH clauses in the original query and multiple-assignment sub-selects in the SET list of UPDATE queries. This is because copying these constructs into a rule query would result in multiple evaluations of the sub-query, contrary to the express intent of the query’s author.

39.4.1. How Update Rules Work

Keep the syntax:

```
CREATE [ OR REPLACE ] RULE name AS ON event
    TO table [ WHERE condition ]
    DO [ ALSO | INSTEAD ] { NOTHING | command | ( command ; command ... ) }
```

in mind. In the following, update rules means rules that are defined on INSERT, UPDATE, or DELETE. Update rules get applied by the rule system when the result relation and the command type of a query tree are equal to the object and event given in the CREATE RULE command. For update rules, the rule system creates a list of query trees. Initially the query-tree list is empty. There can be zero (NOTHING keyword), one, or multiple actions. To simplify, we will look at a rule with one action. This rule can have a qualification or not and it can be INSTEAD or ALSO (the default).

What is a rule qualification? It is a restriction that tells when the actions of the rule should be done and when not. This qualification can only reference the pseudorelations NEW and/or OLD, which basically represent the relation that was given as object (but with a special meaning).

So we have three cases that produce the following query trees for a one-action rule.

No qualification, with either ALSO or INSTEAD

the query tree from the rule action with the original query tree’s qualification added

Qualification given and ALSO

the query tree from the rule action with the rule qualification and the original query tree’s qualification added

Qualification given and INSTEAD

the query tree from the rule action with the rule qualification and the original query tree’s qualification; and the original query tree with the negated rule qualification added
Finally, if the rule is `ALSO`, the unchanged original query tree is added to the list. Since only qualified `INSTEAD` rules already add the original query tree, we end up with either one or two output query trees for a rule with one action.

For `ON INSERT` rules, the original query (if not suppressed by `INSTEAD`) is done before any actions added by rules. This allows the actions to see the inserted row(s). But for `ON UPDATE` and `ON DELETE` rules, the original query is done after the actions added by rules. This ensures that the actions can see the to-be-updated or to-be-deleted rows; otherwise, the actions might do nothing because they find no rows matching their qualifications.

The query trees generated from rule actions are thrown into the rewrite system again, and maybe more rules get applied resulting in more or less query trees. So a rule’s actions must have either a different command type or a different result relation than the rule itself is on, otherwise this recursive process will end up in an infinite loop. (Recursive expansion of a rule will be detected and reported as an error.)

The query trees found in the actions of the `pg_rewrite` system catalog are only templates. Since they can reference the range-table entries for `NEW` and `OLD`, some substitutions have to be made before they can be used. For any reference to `NEW`, the target list of the original query is searched for a corresponding entry. If found, that entry’s expression replaces the reference. Otherwise, `NEW` means the same as `OLD` (for an `UPDATE`) or is replaced by a null value (for an `INSERT`). Any reference to `OLD` is replaced by a reference to the range-table entry that is the result relation.

After the system is done applying update rules, it applies view rules to the produced query tree(s). Views cannot insert new update actions so there is no need to apply update rules to the output of view rewriting.

### 39.4.1.1. A First Rule Step by Step

Say we want to trace changes to the `sl_avail` column in the `shoelace_data` relation. So we set up a log table and a rule that conditionally writes a log entry when an `UPDATE` is performed on `shoelace_data`.

```sql
CREATE TABLE shoelace_log {
    sl_name    text,  -- shoelace changed
    sl_avail   integer,  -- new available value
    log_who    text,  -- who did it
    log_when   timestamp  -- when
};

CREATE RULE log_shoelace AS ON UPDATE TO shoelace_data
    WHERE NEW.sl_avail <> OLD.sl_avail
    DO INSERT INTO shoelace_log VALUES (
        NEW.sl_name,  
        NEW.sl_avail,  
        current_user,  
        current_timestamp
    );
```

Now someone does:

```sql
UPDATE shoelace_data SET sl_avail = 6 WHERE sl_name = 'sl7';
```

and we look at the log table:
That’s what we expected. What happened in the background is the following. The parser created the query tree:

UPDATE shoelace_data SET sl_avail = 6
FROM shoelace_data shoelace_data
WHERE shoelace_data.sl_name = 'sl7';

There is a rule log_shoelace that is ON UPDATE with the rule qualification expression:

NEW.sl_avail <> OLD.sl_avail

and the action:

INSERT INTO shoelace_log VALUES (new.sl_name, new.sl_avail, current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old;

(This looks a little strange since you cannot normally write INSERT ... VALUES ... FROM. The FROM clause here is just to indicate that there are range-table entries in the query tree for new and old. These are needed so that they can be referenced by variables in the INSERT command’s query tree.)

The rule is a qualified ALSO rule, so the rule system has to return two query trees: the modified rule action and the original query tree. In step 1, the range table of the original query is incorporated into the rule’s action query tree. This results in:

INSERT INTO shoelace_log VALUES (new.sl_name, new.sl_avail, current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old, shoelace_data shoelace_data;

In step 2, the rule qualification is added to it, so the result set is restricted to rows where sl_avail changes:

INSERT INTO shoelace_log VALUES (new.sl_name, new.sl_avail, current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old, shoelace_data shoelace_data
WHERE new.sl_avail <> old.sl_avail;

(This looks even stranger, since INSERT ... VALUES doesn’t have a WHERE clause either, but the planner and executor will have no difficulty with it. They need to support this same functionality anyway for INSERT ... SELECT.)

In step 3, the original query tree’s qualification is added, restricting the result set further to only the rows that would have been touched by the original query:

INSERT INTO shoelace_log VALUES (
new.sl_name, new.sl_avail,
current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
shoelace_data shoelace_data
WHERE new.sl_avail <> old.sl_avail
   AND shoelace_data.sl_name = 'sl7';

Step 4 replaces references to NEW by the target list entries from the original query tree or by the
matching variable references from the result relation:

INSERT INTO shoelace_log VALUES (  
   shoelace_data.sl_name, 6,
current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
shoelace_data shoelace_data
WHERE 6 <> old.sl_avail
   AND shoelace_data.sl_name = 'sl7';

Step 5 changes OLD references into result relation references:

INSERT INTO shoelace_log VALUES (  
   shoelace_data.sl_name, 6,
current_user, current_timestamp )
FROM shoelace_data new, shoelace_data old,
shoelace_data shoelace_data
WHERE 6 <> shoelace_data.sl_avail
   AND shoelace_data.sl_name = 'sl7';

That’s it. Since the rule is ALSO, we also output the original query tree. In short, the output from the
rule system is a list of two query trees that correspond to these statements:

INSERT INTO shoelace_log VALUES (  
   shoelace_data.sl_name, 6,
current_user, current_timestamp )
FROM shoelace_data
WHERE 6 <> shoelace_data.sl_avail
   AND shoelace_data.sl_name = 'sl7';

UPDATE shoelace_data SET sl_avail = 6
WHERE sl_name = 'sl7';

These are executed in this order, and that is exactly what the rule was meant to do.
The substitutions and the added qualifications ensure that, if the original query would be, say:

UPDATE shoelace_data SET sl_color = 'green'
WHERE sl_name = 'sl7';

no log entry would get written. In that case, the original query tree does not contain a target list entry
for sl_avail, so NEW.sl_avail will get replaced by shoelace_data.sl_avail. Thus, the extra
command generated by the rule is:

INSERT INTO shoelace_log VALUES (  
   shoelace_data.sl_name, shoelace_data.sl_avail,
and that qualification will never be true. It will also work if the original query modifies multiple rows. So if someone issued the command:

```
UPDATE shoelace_data SET sl_avail = 0
WHERE sl_color = 'black';
```

four rows in fact get updated (sl1, sl2, sl3, and sl4). But sl3 already has sl_avail = 0. In this case, the original query trees qualification is different and that results in the extra query tree:

```
INSERT INTO shoelace_log
SELECT shoelace_data.sl_name, 0,
       current_user, current_timestamp
FROM shoelace_data
WHERE 0 <> shoelace_data.sl_avail
AND shoelace_data.sl_color = 'black';
```

being generated by the rule. This query tree will surely insert three new log entries. And that’s absolutely correct.

Here we can see why it is important that the original query tree is executed last. If the UPDATE had been executed first, all the rows would have already been set to zero, so the logging INSERT would not find any row where 0 <> shoelace_data.sl_avail.

### 39.4.2. Cooperation with Views

A simple way to protect view relations from the mentioned possibility that someone can try to run INSERT, UPDATE, or DELETE on them is to let those query trees get thrown away. So we could create the rules:

```
CREATE RULE shoe_ins_protect AS ON INSERT TO shoe
DO INSTEAD NOTHING;
CREATE RULE shoe_upd_protect AS ON UPDATE TO shoe
DO INSTEAD NOTHING;
CREATE RULE shoe_del_protect AS ON DELETE TO shoe
DO INSTEAD NOTHING;
```

If someone now tries to do any of these operations on the view relation shoe, the rule system will apply these rules. Since the rules have no actions and are INSTEAD, the resulting list of query trees will be empty and the whole query will become nothing because there is nothing left to be optimized or executed after the rule system is done with it.

A more sophisticated way to use the rule system is to create rules that rewrite the query tree into one that does the right operation on the real tables. To do that on the shoelace view, we create the following rules:

```
CREATE RULE shoelace_ins AS ON INSERT TO shoelace
DO INSTEAD
   INSERT INTO shoelace_data VALUES (NEW.sl_name,
                                   NEW.sl_avail,
```

1101
CREATE RULE shoelace_upd AS ON UPDATE TO shoelace
DO INSTEAD
UPDATE shoelace_data
    SET sl_name = NEW.sl_name,
        sl_avail = NEW.sl_avail,
        sl_color = NEW.sl_color,
        sl_len = NEW.sl_len,
        sl_unit = NEW.sl_unit
WHERE sl_name = OLD.sl_name;

CREATE RULE shoelace_del AS ON DELETE TO shoelace
DO INSTEAD
DELETE FROM shoelace_data
WHERE sl_name = OLD.sl_name;

If you want to support RETURNING queries on the view, you need to make the rules include RETURNING clauses that compute the view rows. This is usually pretty trivial for views on a single table, but it’s a bit tedious for join views such as `shoelace`. An example for the insert case is:

CREATE RULE shoelace_ins AS ON INSERT TO shoelace
DO INSTEAD
INSERT INTO shoelace_data VALUES (
    NEW.sl_name,
    NEW.sl_avail,
    NEW.sl_color,
    NEW.sl_len,
    NEW.sl_unit
) RETURNING
    shoelace_data.*,
    (SELECT shoelace_data.sl_len * u.un_fact
     FROM unit u WHERE shoelace_data.sl_unit = u.un_name);

Note that this one rule supports both INSERT and INSERT RETURNING queries on the view — the RETURNING clause is simply ignored for INSERT.

Now assume that once in a while, a pack of shoelaces arrives at the shop and a big parts list along with it. But you don’t want to manually update the `shoelace` view every time. Instead we set up two little tables: one where you can insert the items from the part list, and one with a special trick. The creation commands for these are:

CREATE TABLE shoelace_arrive (
    arr_name text,
    arr_quant integer
);

CREATE TABLE shoelace_ok (
    ok_name text,
    ok_quant integer
);
CREATE RULE shoelace_ok_ins AS ON INSERT TO shoelace_ok
  DO INSTEAD
  UPDATE shoelace
    SET sl_avail = sl_avail + NEW.ok_quant
    WHERE sl_name = NEW.ok_name;

Now you can fill the table shoelace_arrive with the data from the parts list:

SELECT * FROM shoelace_arrive;

<table>
<thead>
<tr>
<th>arr_name</th>
<th>arr_quant</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl3</td>
<td>10</td>
</tr>
<tr>
<td>sl6</td>
<td>20</td>
</tr>
<tr>
<td>sl8</td>
<td>20</td>
</tr>
</tbody>
</table>

(3 rows)

Take a quick look at the current data:

SELECT * FROM shoelace;

<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl1</td>
<td>5</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl2</td>
<td>6</td>
<td>black</td>
<td>100</td>
<td>cm</td>
<td>100</td>
</tr>
<tr>
<td>sl7</td>
<td>6</td>
<td>brown</td>
<td>60</td>
<td>cm</td>
<td>60</td>
</tr>
<tr>
<td>sl3</td>
<td>0</td>
<td>black</td>
<td>35</td>
<td>inch</td>
<td>88.9</td>
</tr>
<tr>
<td>sl4</td>
<td>8</td>
<td>black</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl8</td>
<td>1</td>
<td>brown</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl5</td>
<td>4</td>
<td>brown</td>
<td>1</td>
<td>m</td>
<td>100</td>
</tr>
<tr>
<td>sl6</td>
<td>0</td>
<td>brown</td>
<td>0.9</td>
<td>m</td>
<td>90</td>
</tr>
</tbody>
</table>

(8 rows)

Now move the arrived shoelaces in:

INSERT INTO shoelace_ok SELECT * FROM shoelace_arrive;

and check the results:

SELECT * FROM shoelace ORDER BY sl_name;

<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl1</td>
<td>5</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl2</td>
<td>6</td>
<td>black</td>
<td>100</td>
<td>cm</td>
<td>100</td>
</tr>
<tr>
<td>sl7</td>
<td>6</td>
<td>brown</td>
<td>60</td>
<td>cm</td>
<td>60</td>
</tr>
<tr>
<td>sl4</td>
<td>8</td>
<td>black</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl3</td>
<td>10</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl8</td>
<td>21</td>
<td>brown</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
<tr>
<td>sl5</td>
<td>4</td>
<td>brown</td>
<td>1</td>
<td>m</td>
<td>100</td>
</tr>
<tr>
<td>sl6</td>
<td>20</td>
<td>brown</td>
<td>0.9</td>
<td>m</td>
<td>90</td>
</tr>
</tbody>
</table>

(8 rows)

SELECT * FROM shoelace_log;

<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>log_who</th>
<th>log_when</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl7</td>
<td>6</td>
<td>Al</td>
<td>Tue Oct 20 19:14:45 1998 MET DST</td>
</tr>
</tbody>
</table>
It’s a long way from the one \texttt{INSERT ... SELECT} to these results. And the description of the query-tree transformation will be the last in this chapter. First, there is the parser’s output:

\texttt{INSERT INTO shoelace\_ok}
\texttt{SELECT shoelace\_arrive.arr\_name, shoelace\_arrive.arr\_quant}
\texttt{FROM shoelace\_arrive shoelace\_arrive, shoelace\_ok shoelace\_ok;}

Now the first rule \texttt{shoelace\_ok\_ins} is applied and turns this into:

\texttt{UPDATE shoelace}
\texttt{SET sl\_avail = shoelace.sl\_avail + shoelace\_arrive.arr\_quant}
\texttt{FROM shoelace\_arrive shoelace\_arrive, shoelace\_ok shoelace\_ok,}
\texttt{shoelace\_ok old, shoelace\_ok new,}
\texttt{shoelace shoelace}
\texttt{WHERE shoelace.sl\_name = shoelace\_arrive.arr\_name;}

and throws away the original \texttt{INSERT on shoelace\_ok}. This rewritten query is passed to the rule system again, and the second applied rule \texttt{shoelace\_upd} produces:

\texttt{UPDATE shoelace\_data}
\texttt{SET sl\_name = shoelace.sl\_name,}
\texttt{sl\_avail = shoelace.sl\_avail + shoelace\_arrive.arr\_quant,}
\texttt{sl\_color = shoelace.sl\_color,}
\texttt{sl\_len = shoelace.sl\_len,}
\texttt{sl\_unit = shoelace.sl\_unit}
\texttt{FROM shoelace\_arrive shoelace\_arrive, shoelace\_ok shoelace\_ok,}
\texttt{shoelace\_ok old, shoelace\_ok new,}
\texttt{shoelace shoelace, shoelace old,}
\texttt{shoelace new, shoelace\_data shoelace\_data}
\texttt{WHERE shoelace.sl\_name = shoelace\_arrive.arr\_name}
\texttt{AND shoelace\_data.sl\_name = shoelace\_ok\_ins;}

Again it’s an \texttt{INSTEAD} rule and the previous query tree is trashed. Note that this query still uses the view \texttt{shoelace}. But the rule system isn’t finished with this step, so it continues and applies the \texttt{_RETURN} rule on it, and we get:

\texttt{UPDATE shoelace\_data}
\texttt{SET sl\_name = s.sl\_name,}
\texttt{sl\_avail = s.sl\_avail + shoelace\_arrive.arr\_quant,}
\texttt{sl\_color = s.sl\_color,}
\texttt{sl\_len = s.sl\_len,}
\texttt{sl\_unit = s.sl\_unit}
\texttt{FROM shoelace\_arrive shoelace\_arrive, shoelace\_ok shoelace\_ok,}
\texttt{shoelace\_ok old, shoelace\_ok new,}
\texttt{shoelace shoelace, shoelace old,}
\texttt{shoelace new, shoelace\_data shoelace\_data,}
\texttt{shoelace old, shoelace new,}
\texttt{shoelace\_data s, unit u}
\texttt{WHERE s.sl\_name = shoelace\_arrive.arr\_name}
\texttt{AND shoelace\_data.sl\_name = s.sl\_name;}

1104
Chapter 39. The Rule System

Finally, the rule `log_shoelace` gets applied, producing the extra query tree:

```sql
INSERT INTO shoelace_log
SELECT s.sl_name,
       s.sl_avail + shoelace_arrive.arr_quant,
       current_user,
       current_timestamp
FROM shoelace_arrive shoelace_arrive, shoelace_data shoelace_data,
     shoelace ok shoelace ok, shoelace ok old, shoelace ok new,
     shoelace new, shoelace_data shoelace_data, shoelace old, shoelace new,
     shoelace_data s, unit u,
     shoelace_data old, shoelace_data new
WHERE s.sl_name = shoelace_arrive.arr_name
AND shoelace_data.sl_name = s.sl_name
AND (s.sl_avail + shoelace_arrive.arr_quant) <> s.sl_avail;
```

After that the rule system runs out of rules and returns the generated query trees.

So we end up with two final query trees that are equivalent to the SQL statements:

```sql
UPDATE shoelace_data
SET sl_avail = shoelace_data.sl_avail + shoelace_arrive.arr_quant
FROM shoelace_arrive shoelace_arrive,
     shoelace_data shoelace_data,
     shoelace_data s
WHERE s.sl_name = shoelace_arrive.arr_name
AND shoelace_data.sl_name = s.sl_name
AND shoelace_arrive.sl_name = s.sl_name;
```

The result is that data coming from one relation inserted into another, changed into updates on a third, changed into updating a fourth plus logging that final update in a fifth gets reduced into two queries.

There is a little detail that’s a bit ugly. Looking at the two queries, it turns out that the `shoelace_data` relation appears twice in the range table where it could definitely be reduced to one. The planner does not handle it and so the execution plan for the rule systems output of the `INSERT` will be

```
Nested Loop
  -> Merge Join
    -> Seq Scan
      -> Sort
    -> Seq Scan on s
  -> Seq Scan
    -> Sort
      -> Seq Scan on shoelace_arrive
  -> Seq Scan on shoelace_data
```
while omitting the extra range table entry would result in a

Merge Join
-\> Seq Scan
-\> Sort
-\> Seq Scan on s
-\> Seq Scan
-\> Sort
-\> Seq Scan on shoelace_arrive

which produces exactly the same entries in the log table. Thus, the rule system caused one extra scan on the table shoelace_data that is absolutely not necessary. And the same redundant scan is done once more in the UPDATE. But it was a really hard job to make that all possible at all.

Now we make a final demonstration of the PostgreSQL rule system and its power. Say you add some shoelaces with extraordinary colors to your database:

```
INSERT INTO shoelace VALUES ('sl9', 0, 'pink', 35.0, 'inch', 0.0);
INSERT INTO shoelace VALUES ('sl10', 1000, 'magenta', 40.0, 'inch', 0.0);
```

We would like to make a view to check which shoelace entries do not fit any shoe in color. The view for this is:

```
CREATE VIEW shoelace_mismatch AS
    SELECT * FROM shoelace WHERE NOT EXISTS
        (SELECT shoename FROM shoe WHERE slcolor = sl_color);
```

Its output is:

```
SELECT * FROM shoelace_mismatch;

<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl9</td>
<td>0</td>
<td>pink</td>
<td>35</td>
<td>inch</td>
<td>88.9</td>
</tr>
<tr>
<td>sl10</td>
<td>1000</td>
<td>magenta</td>
<td>40</td>
<td>inch</td>
<td>101.6</td>
</tr>
</tbody>
</table>
```

Now we want to set it up so that mismatching shoelaces that are not in stock are deleted from the database. To make it a little harder for PostgreSQL, we don’t delete it directly. Instead we create one more view:

```
CREATE VIEW shoelace_can_delete AS
    SELECT * FROM shoelace_mismatch WHERE sl_avail = 0;
```

and do it this way:

```
DELETE FROM shoelace WHERE EXISTS
    (SELECT * FROM shoelace_can_delete
     WHERE sl_name = shoelace.sl_name);
```

Voilà:

```
SELECT * FROM shoelace;

<table>
<thead>
<tr>
<th>sl_name</th>
<th>sl_avail</th>
<th>sl_color</th>
<th>sl_len</th>
<th>sl_unit</th>
<th>sl_len_cm</th>
</tr>
</thead>
<tbody>
<tr>
<td>sl1</td>
<td>5</td>
<td>black</td>
<td>80</td>
<td>cm</td>
<td>80</td>
</tr>
<tr>
<td>sl2</td>
<td>6</td>
<td>black</td>
<td>100</td>
<td>cm</td>
<td>100</td>
</tr>
</tbody>
</table>
```

1106
A DELETE on a view, with a subquery qualification that in total uses 4 nesting/joined views, where one of them itself has a subquery qualification containing a view and where calculated view columns are used, gets rewritten into one single query tree that deletes the requested data from a real table.

There are probably only a few situations out in the real world where such a construct is necessary. But it makes you feel comfortable that it works.

### 39.5. Rules and Privileges

Due to rewriting of queries by the PostgreSQL rule system, other tables/views than those used in the original query get accessed. When update rules are used, this can include write access to tables.

Rewrite rules don’t have a separate owner. The owner of a relation (table or view) is automatically the owner of the rewrite rules that are defined for it. The PostgreSQL rule system changes the behavior of the default access control system. Relations that are used due to rules get checked against the privileges of the rule owner, not the user invoking the rule. This means that users only need the required privileges for the tables/views that are explicitly named in their queries.

For example: A user has a list of phone numbers where some of them are private, the others are of interest for the assistant of the office. The user can construct the following:

```sql
CREATE TABLE phone_data (person text, phone text, private boolean);
CREATE VIEW phone_number AS
    SELECT person, CASE WHEN NOT private THEN phone END AS phone
    FROM phone_data;
GRANT SELECT ON phone_number TO assistant;
```

Nobody except that user (and the database superusers) can access the `phone_data` table. But because of the GRANT, the assistant can run a SELECT on the `phone_number` view. The rule system will rewrite the SELECT from `phone_number` into a SELECT from `phone_data`. Since the user is the owner of `phone_number` and therefore the owner of the rule, the read access to `phone_data` is now checked against the user’s privileges and the query is permitted. The check for accessing `phone_number` is also performed, but this is done against the invoking user, so nobody but the user and the assistant can use it.

The privileges are checked rule by rule. So the assistant is for now the only one who can see the public phone numbers. But the assistant can set up another view and grant access to that to the public. Then, anyone can see the `phone_number` data through the assistant’s view. What the assistant cannot do is to create a view that directly accesses `phone_data`. (Actually the assistant can, but it will not work since every access will be denied during the permission checks.) And as soon as the user notices that the assistant opened their `phone_number` view, the user can revoke the assistant’s access. Immediately, any access to the assistant’s view would fail.
Chapter 39. The Rule System

One might think that this rule-by-rule checking is a security hole, but in fact it isn’t. But if it did not work this way, the assistant could set up a table with the same columns as `phone_number` and copy the data to there once per day. Then it’s the assistant’s own data and the assistant can grant access to everyone they want. A `GRANT` command means, “I trust you”. If someone you trust does the thing above, it’s time to think it over and then use `REVOKE`.

Note that while views can be used to hide the contents of certain columns using the technique shown above, they cannot be used to reliably conceal the data in unseen rows unless the `security_barrier` flag has been set. For example, the following view is insecure:

```sql
CREATE VIEW phone_number AS
    SELECT person, phone FROM phone_data WHERE phone NOT LIKE '412%';
```

This view might seem secure, since the rule system will rewrite any `SELECT` from `phone_number` into a `SELECT` from `phone_data` and add the qualification that only entries where `phone` does not begin with 412 are wanted. But if the user can create their own functions, it is not difficult to convince the planner to execute the user-defined function prior to the `NOT LIKE` expression. For example:

```sql
CREATE FUNCTION tricky(text, text) RETURNS bool AS $$
BEGIN
    RAISE NOTICE '% = %', $1, $2;
    RETURN true;
END
$$ LANGUAGE plpgsql COST 0.0000000000000000000001;

SELECT * FROM phone_number WHERE tricky(person, phone);
```

Every person and phone number in the `phone_data` table will be printed as a `NOTICE`, because the planner will choose to execute the inexpensive `tricky` function before the more expensive `NOT LIKE`. Even if the user is prevented from defining new functions, built-in functions can be used in similar attacks. (For example, most casting functions include their input values in the error messages they produce.)

Similar considerations apply to update rules. In the examples of the previous section, the owner of the tables in the example database could grant the privileges `SELECT`, `INSERT`, `UPDATE`, and `DELETE` on the `shoelace` view to someone else, but only `SELECT` on `shoelace_log`. The rule action to write log entries will still be executed successfully, and that other user could see the log entries. But they could not create fake entries, nor could they manipulate or remove existing ones. In this case, there is no possibility of subverting the rules by convincing the planner to alter the order of operations, because the only rule which references `shoelace_log` is an unqualified `INSERT`. This might not be true in more complex scenarios.

When it is necessary for a view to provide row level security, the `security_barrier` attribute should be applied to the view. This prevents maliciously-chosen functions and operators from being passed values from rows until after the view has done its work. For example, if the view shown above had been created like this, it would be secure:

```sql
CREATE VIEW phone_number WITH (security_barrier) AS
    SELECT person, phone FROM phone_data WHERE phone NOT LIKE '412%';
```

Views created with the `security_barrier` may perform far worse than views created without this option. In general, there is no way to avoid this: the fastest possible plan must be rejected if it may compromise security. For this reason, this option is not enabled by default.

The query planner has more flexibility when dealing with functions that have no side effects. Such functions are referred to as `LEAKPROOF`, and include many simple, commonly used operators, such as many equality operators. The query planner can safely allow such functions to be evaluated at any
point in the query execution process, since invoking them on rows invisible to the user will not leak any information about the unseen rows. Further, functions which do not take arguments or which are not passed any arguments from the security barrier view do not have to be marked as LEAKPROOF to be pushed down, as they never receive data from the view. In contrast, a function that might throw an error depending on the values received as arguments (such as one that throws an error in the event of overflow or division by zero) is not leak-proof, and could provide significant information about the unseen rows if applied before the security view’s row filters.

It is important to understand that even a view created with the security_barrier option is intended to be secure only in the limited sense that the contents of the invisible tuples will not be passed to possibly-insecure functions. The user may well have other means of making inferences about the unseen data; for example, they can see the query plan using EXPLAIN, or measure the run time of queries against the view. A malicious attacker might be able to infer something about the amount of unseen data, or even gain some information about the data distribution or most common values (since these things may affect the run time of the plan; or even, since they are also reflected in the optimizer statistics, the choice of plan). If these types of "covert channel" attacks are of concern, it is probably unwise to grant any access to the data at all.

39.6. Rules and Command Status

The PostgreSQL server returns a command status string, such as INSERT 149592 1, for each command it receives. This is simple enough when there are no rules involved, but what happens when the query is rewritten by rules?

Rules affect the command status as follows:

- If there is no unconditional INSTEAD rule for the query, then the originally given query will be executed, and its command status will be returned as usual. (But note that if there were any conditional INSTEAD rules, the negation of their qualifications will have been added to the original query. This might reduce the number of rows it processes, and if so the reported status will be affected.)

- If there is any unconditional INSTEAD rule for the query, then the original query will not be executed at all. In this case, the server will return the command status for the last query that was inserted by an INSTEAD rule (conditional or unconditional) and is of the same command type (INSERT, UPDATE, or DELETE) as the original query. If no query meeting those requirements is added by any rule, then the returned command status shows the original query type and zeroes for the row-count and OID fields.

The programmer can ensure that any desired INSTEAD rule is the one that sets the command status in the second case, by giving it the alphabetically last rule name among the active rules, so that it gets applied last.

39.7. Rules Versus Triggers

Many things that can be done using triggers can also be implemented using the PostgreSQL rule system. One of the things that cannot be implemented by rules are some kinds of constraints, especially foreign keys. It is possible to place a qualified rule that rewrites a command to NOTHING if the value of a column does not appear in another table. But then the data is silently thrown away and that’s not a
good idea. If checks for valid values are required, and in the case of an invalid value an error message should be generated, it must be done by a trigger.

In this chapter, we focused on using rules to update views. All of the update rule examples in this chapter can also be implemented using INSTEAD OF triggers on the views. Writing such triggers is often easier than writing rules, particularly if complex logic is required to perform the update.

For the things that can be implemented by both, which is best depends on the usage of the database. A trigger is fired once for each affected row. A rule modifies the query or generates an additional query. So if many rows are affected in one statement, a rule issuing one extra command is likely to be faster than a trigger that is called for every single row and must re-determine what to do many times. However, the trigger approach is conceptually far simpler than the rule approach, and is easier for novices to get right.

Here we show an example of how the choice of rules versus triggers plays out in one situation. There are two tables:

```
CREATE TABLE computer {
    hostname text, -- indexed
    manufacturer text -- indexed
};
CREATE TABLE software {
    software text, -- indexed
    hostname text -- indexed
};
```

Both tables have many thousands of rows and the indexes on `hostname` are unique. The rule or trigger should implement a constraint that deletes rows from `software` that reference a deleted computer.

The trigger would use this command:

```
DELETE FROM software WHERE hostname = $1;
```

Since the trigger is called for each individual row deleted from `computer`, it can prepare and save the plan for this command and pass the `hostname` value in the parameter. The rule would be written as:

```
CREATE RULE computer_del AS ON DELETE TO computer
    DO DELETE FROM software WHERE hostname = OLD.hostname;
```

Now we look at different types of deletes. In the case of a:

```
DELETE FROM computer WHERE hostname = ‘mypc.local.net’;
```

the table `computer` is scanned by index (fast), and the command issued by the trigger would also use an index scan (also fast). The extra command from the rule would be:

```
DELETE FROM software WHERE computer.hostname = ‘mypc.local.net’
    AND software.hostname = computer.hostname;
```

Since there are appropriate indexes set up, the planner will create a plan of

```
Nestloop
  --> Index Scan using comp_hostidx on computer
  --> Index Scan using soft_hostidx on software
```

So there would be not that much difference in speed between the trigger and the rule implementation.
With the next delete we want to get rid of all the 2000 computers where the `hostname` starts with `old`. There are two possible commands to do that. One is:

```sql
DELETE FROM computer WHERE hostname >= 'old'
  AND hostname < 'ole'
```

The command added by the rule will be:

```sql
DELETE FROM software WHERE computer.hostname >= 'old' AND computer.hostname < 'ole'
  AND software.hostname = computer.hostname;
```

with the plan:

```
Hash Join
  ->  Seq Scan on software
  ->  Hash
      ->  Index Scan using comp_hostidx on computer
```

The other possible command is:

```sql
DELETE FROM computer WHERE hostname ~ '^old';
```

which results in the following executing plan for the command added by the rule:

```
Nestloop
  ->  Index Scan using comp_hostidx on computer
  ->  Index Scan using soft_hostidx on software
```

This shows, that the planner does not realize that the qualification for `hostname` in `computer` could also be used for an index scan on `software` when there are multiple qualification expressions combined with `AND`, which is what it does in the regular-expression version of the command. The trigger will get invoked once for each of the 2000 old computers that have to be deleted, and that will result in one index scan over `computer` and 2000 index scans over `software`. The rule implementation will do it with two commands that use indexes. And it depends on the overall size of the table `software` whether the rule will still be faster in the sequential scan situation. 2000 command executions from the trigger over the SPI manager take some time, even if all the index blocks will soon be in the cache.

The last command we look at is:

```sql
DELETE FROM computer WHERE manufacturer = 'bim';
```

Again this could result in many rows to be deleted from `computer`. So the trigger will again run many commands through the executor. The command generated by the rule will be:

```sql
DELETE FROM software WHERE computer.manufacturer = 'bim'
  AND software.hostname = computer.hostname;
```

The plan for that command will again be the nested loop over two index scans, only using a different index on `computer`:

```
Nestloop
  ->  Index Scan using comp_manufidx on computer
  ->  Index Scan using soft_hostidx on software
```

In any of these cases, the extra commands from the rule system will be more or less independent from the number of affected rows in a command.
The summary is, rules will only be significantly slower than triggers if their actions result in large and badly qualified joins, a situation where the planner fails.
Chapter 40. Procedural Languages

PostgreSQL allows user-defined functions to be written in other languages besides SQL and C. These other languages are generically called procedural languages (PLs). For a function written in a procedural language, the database server has no built-in knowledge about how to interpret the function’s source text. Instead, the task is passed to a special handler that knows the details of the language. The handler could either do all the work of parsing, syntax analysis, execution, etc. itself, or it could serve as “glue” between PostgreSQL and an existing implementation of a programming language. The handler itself is a C language function compiled into a shared object and loaded on demand, just like any other C function.

There are currently four procedural languages available in the standard PostgreSQL distribution: PL/pgSQL (Chapter 41), PL/Tcl (Chapter 42), PL/Perl (Chapter 43), and PL/Python (Chapter 44). There are additional procedural languages available that are not included in the core distribution. Appendix H has information about finding them. In addition other languages can be defined by users; the basics of developing a new procedural language are covered in Chapter 54.

40.1. Installing Procedural Languages

A procedural language must be “installed” into each database where it is to be used. But procedural languages installed in the database template1 are automatically available in all subsequently created databases, since their entries in template1 will be copied by CREATE DATABASE. So the database administrator can decide which languages are available in which databases and can make some languages available by default if desired.

For the languages supplied with the standard distribution, it is only necessary to execute CREATE EXTENSION language_name to install the language into the current database. Alternatively, the program createlang can be used to do this from the shell command line. For example, to install the language PL/Perl into the database template1, use:

createlang plperl template1

The manual procedure described below is only recommended for installing languages that have not been packaged as extensions.

Manual Procedural Language Installation

A procedural language is installed in a database in five steps, which must be carried out by a database superuser. In most cases the required SQL commands should be packaged as the installation script of an “extension”, so that CREATE EXTENSION can be used to execute them.

1. The shared object for the language handler must be compiled and installed into an appropriate library directory. This works in the same way as building and installing modules with regular user-defined C functions does; see Section 36.9.6. Often, the language handler will depend on an external library that provides the actual programming language engine; if so, that must be installed as well.

2. The handler must be declared with the command

```sql
CREATE FUNCTION handler_function_name()
RETURNS language_handler
AS 'path-to-shared-object'
LANGUAGE C;
```
The special return type of `language_handler` tells the database system that this function does not return one of the defined SQL data types and is not directly usable in SQL statements.

3. Optionally, the language handler can provide an “inline” handler function that executes anonymous code blocks (DO commands) written in this language. If an inline handler function is provided by the language, declare it with a command like

```sql
CREATE FUNCTION inline_function_name(internal)
    RETURNS void
    AS 'path-to-shared-object'
    LANGUAGE C;
```

4. Optionally, the language handler can provide a “validator” function that checks a function definition for correctness without actually executing it. The validator function is called by `CREATE FUNCTION` if it exists. If a validator function is provided by the language, declare it with a command like

```sql
CREATE FUNCTION validator_function_name(oid)
    RETURNS void
    AS 'path-to-shared-object'
    LANGUAGE C STRICT;
```

5. Finally, the PL must be declared with the command

```sql
CREATE [TRUSTED] [PROCEDURAL] LANGUAGE language_name
    HANDLER handler_function_name
    [INLINE inline_function_name]
    [VALIDATOR validator_function_name];
```

The optional key word `TRUSTED` specifies that the language does not grant access to data that the user would not otherwise have. Trusted languages are designed for ordinary database users (those without superuser privilege) and allows them to safely create functions and trigger procedures. Since PL functions are executed inside the database server, the `TRUSTED` flag should only be given for languages that do not allow access to database server internals or the file system. The languages PL/pgSQL, PL/Tcl, and PL/Perl are considered trusted; the languages PL/TclU, PL/PerlU, and PL/PythonU are designed to provide unlimited functionality and should not be marked trusted.

Example 40-1 shows how the manual installation procedure would work with the language PL/Perl.

**Example 40-1. Manual Installation of PL/Perl**

The following command tells the database server where to find the shared object for the PL/Perl language’s call handler function:

```sql
CREATE FUNCTION plperl_call_handler() RETURNS language_handler AS
    '$libdir/plperl' LANGUAGE C;
```

PL/Perl has an inline handler function and a validator function, so we declare those too:

```sql
CREATE FUNCTION plperl_inline_handler(internal) RETURNS void AS
    '$libdir/plperl' LANGUAGE C;

CREATE FUNCTION plperl_validator(oid) RETURNS void AS
    '$libdir/plperl' LANGUAGE C STRICT;
```

The command:

```sql
CREATE TRUSTED PROCEDURAL LANGUAGE plperl
    HANDLER plperl_call_handler
    INLINE plperl_inline_handler
    VALIDATOR plperl_validator;
```
then defines that the previously declared functions should be invoked for functions and trigger procedures where the language attribute is `plperl`.

In a default PostgreSQL installation, the handler for the PL/pgSQL language is built and installed into the “library” directory; furthermore, the PL/pgSQL language itself is installed in all databases. If Tcl support is configured in, the handlers for PL/Tcl and PL/TcU are built and installed in the library directory, but the language itself is not installed in any database by default. Likewise, the PL/Perl and PL/PerlU handlers are built and installed if Perl support is configured, and the PL/PythonU handler is installed if Python support is configured, but these languages are not installed by default.
Chapter 41. PL/pgSQL - SQL Procedural Language

41.1. Overview

PL/pgSQL is a loadable procedural language for the PostgreSQL database system. The design goals of PL/pgSQL were to create a loadable procedural language that

• can be used to create functions and trigger procedures,
• adds control structures to the SQL language,
• can perform complex computations,
• inherits all user-defined types, functions, and operators,
• can be defined to be trusted by the server,
• is easy to use.

Functions created with PL/pgSQL can be used anywhere that built-in functions could be used. For example, it is possible to create complex conditional computation functions and later use them to define operators or use them in index expressions.

In PostgreSQL 9.0 and later, PL/pgSQL is installed by default. However it is still a loadable module, so especially security-conscious administrators could choose to remove it.

41.1.1. Advantages of Using PL/pgSQL

SQL is the language PostgreSQL and most other relational databases use as query language. It’s portable and easy to learn. But every SQL statement must be executed individually by the database server.

That means that your client application must send each query to the database server, wait for it to be processed, receive and process the results, do some computation, then send further queries to the server. All this incurs interprocess communication and will also incur network overhead if your client is on a different machine than the database server.

With PL/pgSQL, you can group a block of computation and a series of queries inside the database server, thus having the power of a procedural language and the ease of use of SQL, but with considerable savings of client/server communication overhead.

• Extra round trips between client and server are eliminated
• Intermediate results that the client does not need do not have to be marshaled or transferred between server and client
• Multiple rounds of query parsing can be avoided

This can result in a considerable performance increase as compared to an application that does not use stored functions.

Also, with PL/pgSQL you can use all the data types, operators and functions of SQL.
41.1.2. Supported Argument and Result Data Types

Functions written in PL/pgSQL can accept as arguments any scalar or array data type supported by the server, and they can return a result of any of these types. They can also accept or return any composite type (row type) specified by name. It is also possible to declare a PL/pgSQL function as returning record, which means that the result is a row type whose columns are determined by specification in the calling query, as discussed in Section 7.2.1.4.

PL/pgSQL functions can be declared to accept a variable number of arguments by using the VARIADIC marker. This works exactly the same way as for SQL functions, as discussed in Section 36.4.5.

PL/pgSQL functions can also be declared to accept and return the polymorphic types anyelement, anyarray, anynonarray, anyenum, and anyrange. The actual data types handled by a polymorphic function can vary from call to call, as discussed in Section 36.2.5. An example is shown in Section 41.3.1.

PL/pgSQL functions can also be declared to return a “set” (or table) of any data type that can be returned as a single instance. Such a function generates its output by executing RETURN NEXT for each desired element of the result set, or by using RETURN QUERY to output the result of evaluating a query.

Finally, a PL/pgSQL function can be declared to return void if it has no useful return value.

PL/pgSQL functions can also be declared with output parameters in place of an explicit specification of the return type. This does not add any fundamental capability to the language, but it is often convenient, especially for returning multiple values. The RETURNS TABLE notation can also be used in place of RETURNS SETOF.

Specific examples appear in Section 41.3.1 and Section 41.6.1.

41.2. Structure of PL/pgSQL

Functions written in PL/pgSQL are defined to the server by executing CREATE FUNCTION commands. Such a command would normally look like, say,

```sql
CREATE FUNCTION somefunc(integer, text) RETURNS integer
AS 'function body text'
LANGUAGE plpgsql;
```

The function body is simply a string literal so far as CREATE FUNCTION is concerned. It is often helpful to use dollar quoting (see Section 4.1.2.4) to write the function body, rather than the normal single quote syntax. Without dollar quoting, any single quotes or backslashes in the function body must be escaped by doubling them. Almost all the examples in this chapter use dollar-quoted literals for their function bodies.

PL/pgSQL is a block-structured language. The complete text of a function body must be a block. A block is defined as:

```sql
[ <<label>> ]
[ DECLARE
   declarations ]
BEGIN
   statements
END [ label ];
```
Each declaration and each statement within a block is terminated by a semicolon. A block that appears within another block must have a semicolon after `END`, as shown above; however the final `END` that concludes a function body does not require a semicolon.

**Tip:** A common mistake is to write a semicolon immediately after `BEGIN`. This is incorrect and will result in a syntax error.

A label is only needed if you want to identify the block for use in an `EXIT` statement, or to qualify the names of the variables declared in the block. If a label is given after `END`, it must match the label at the block’s beginning.

All key words are case-insensitive. Identifiers are implicitly converted to lower case unless double-quoted, just as they are in ordinary SQL commands.

Comments work the same way in PL/pgSQL code as in ordinary SQL. A double dash (`--`) starts a comment that extends to the end of the line. A `/*` starts a block comment that extends to the matching occurrence of `*/`. Block comments nest.

Any statement in the statement section of a block can be a subblock. Subblocks can be used for logical grouping or to localize variables to a small group of statements. Variables declared in a subblock mask any similarly-named variables of outer blocks for the duration of the subblock; but you can access the outer variables anyway if you qualify their names with their block’s label. For example:

```sql
CREATE FUNCTION somefunc() RETURNS integer AS $$
<< outerblock >>
DECLARE
textarea integer := 30;
BEGIN
  RAISE NOTICE 'TextArea here is %', textarea; -- Prints 30
  textarea := 50;
  -- Create a subblock
  --
  DECLARE
textarea integer := 80;
  BEGIN
    RAISE NOTICE 'TextArea here is %', textarea; -- Prints 80
    RAISE NOTICE 'Outer textarea here is %', outerblock.textarea; -- Prints 50
  END;

  RAISE NOTICE 'TextArea here is %', textarea; -- Prints 50

  RETURN textarea;
END;
$$ LANGUAGE plpgsql;
```

**Note:** There is actually a hidden “outer block” surrounding the body of any PL/pgSQL function. This block provides the declarations of the function’s parameters (if any), as well as some special variables such as `FOUND` (see Section 41.5.5). The outer block is labeled with the function’s name, meaning that parameters and special variables can be qualified with the function’s name.
It is important not to confuse the use of `BEGIN/END` for grouping statements in PL/pgSQL with the similarly-named SQL commands for transaction control. PL/pgSQL’s `BEGIN/END` are only for grouping; they do not start or end a transaction. Functions and trigger procedures are always executed within a transaction established by an outer query — they cannot start or commit that transaction, since there would be no context for them to execute in. However, a block containing an `EXCEPTION` clause effectively forms a subtransaction that can be rolled back without affecting the outer transaction. For more about that see Section 41.6.6.

### 41.3. Declarations

All variables used in a block must be declared in the declarations section of the block. (The only exceptions are that the loop variable of a `FOR` loop iterating over a range of integer values is automatically declared as an integer variable, and likewise the loop variable of a `FOR` loop iterating over a cursor’s result is automatically declared as a record variable.)

PL/pgSQL variables can have any SQL data type, such as `integer`, `varchar`, and `char`.

Here are some examples of variable declarations:

```sql
user_id integer;
quantity numeric(5);
url varchar;
myrow tablename%ROWTYPE;
myfield tablename.columnname%TYPE;
arow RECORD;
```

The general syntax of a variable declaration is:

```sql
name [ CONSTANT ] type [ COLLATE collation_name ] [ NOT NULL ] [ { DEFAULT | := | = } expression ];
```

The `DEFAULT` clause, if given, specifies the initial value assigned to the variable when the block is entered. If the `DEFAULT` clause is not given then the variable is initialized to the SQL null value. The `CONSTANT` option prevents the variable from being assigned to after initialization, so that its value will remain constant for the duration of the block. The `COLLATE` option specifies a collation to use for the variable (see Section 41.3.6). If `NOT NULL` is specified, an assignment of a null value results in a run-time error. All variables declared as `NOT NULL` must have a nonnull default value specified. Equal (`=`) can be used instead of PL/SQL-compliant `:=`.

A variable’s default value is evaluated and assigned to the variable each time the block is entered (not just once per function call). So, for example, assigning `now()` to a variable of type `timestamp` causes the variable to have the time of the current function call, not the time when the function was precompiled.

Examples:

```sql
quantity integer DEFAULT 32;
url varchar := 'http://mysite.com';
user_id CONSTANT integer := 10;
```
41.3.1. Declaring Function Parameters

Parameters passed to functions are named with the identifiers $1, $2, etc. Optionally, aliases can be declared for $n parameter names for increased readability. Either the alias or the numeric identifier can then be used to refer to the parameter value.

There are two ways to create an alias. The preferred way is to give a name to the parameter in the CREATE FUNCTION command, for example:

```sql
CREATE FUNCTION sales_tax(subtotal real) RETURNS real AS $$
BEGIN
    RETURN subtotal * 0.06;
END;
$$ LANGUAGE plpgsql;
```

The other way is to explicitly declare an alias, using the declaration syntax

```sql
name ALIAS FOR $n;
```

The same example in this style looks like:

```sql
CREATE FUNCTION sales_tax(real) RETURNS real AS $$
DECLARE
    subtotal ALIAS FOR $1;
BEGIN
    RETURN subtotal * 0.06;
END;
$$ LANGUAGE plpgsql;
```

Note: These two examples are not perfectly equivalent. In the first case, subtotal could be referenced as sales_tax.subtotal, but in the second case it could not. (Had we attached a label to the inner block, subtotal could be qualified with that label, instead.)

Some more examples:

```sql
CREATE FUNCTION instr(varchar, integer) RETURNS integer AS $$
DECLARE
    v_string ALIAS FOR $1;
    index ALIAS FOR $2;
BEGIN
    -- some computations using v_string and index here
END;
$$ LANGUAGE plpgsql;
```

```sql
CREATE FUNCTION concat_selected_fields(in_t sometablename) RETURNS text AS $$
BEGIN
    RETURN in_t.f1 || in_t.f3 || in_t.f5 || in_t.f7;
END;
$$ LANGUAGE plpgsql;
```

When a PL/pgSQL function is declared with output parameters, the output parameters are given $n names and optional aliases in just the same way as the normal input parameters. An output parameter
is effectively a variable that starts out NULL; it should be assigned to during the execution of the function. The final value of the parameter is what is returned. For instance, the sales-tax example could also be done this way:

```sql
CREATE FUNCTION sales_tax(subtotal real, OUT tax real) AS $$
BEGIN
  tax := subtotal * 0.06;
END;
$$ LANGUAGE plpgsql;
```

Notice that we omitted `RETURNS real` — we could have included it, but it would be redundant.

Output parameters are most useful when returning multiple values. A trivial example is:

```sql
CREATE FUNCTION sum_n_product(x int, y int, OUT sum int, OUT prod int) AS $$
BEGIN
  sum := x + y;
  prod := x * y;
END;
$$ LANGUAGE plpgsql;
```

As discussed in Section 36.4.4, this effectively creates an anonymous record type for the function’s results. If a `RETURNS` clause is given, it must say `RETURNS record`.

Another way to declare a PL/pgSQL function is with `RETURNS TABLE`, for example:

```sql
CREATE FUNCTION extended_sales(p_itemno int)
RETURNS TABLE(quantity int, total numeric) AS $$
BEGIN
  RETURN QUERY SELECT s.quantity, s.quantity * s.price FROM sales AS s
  WHERE s.itemno = p_itemno;
END;
$$ LANGUAGE plpgsql;
```

This is exactly equivalent to declaring one or more `OUT` parameters and specifying `RETURNS SETOF sometype`.

When the return type of a PL/pgSQL function is declared as a polymorphic type (`anyelement`, `anyarray`, `anynonarray`, `anyenum`, or `anyrange`), a special parameter `$0` is created. Its data type is the actual return type of the function, as deduced from the actual input types (see Section 36.2.5). This allows the function to access its actual return type as shown in Section 41.3.3. `$0` is initialized to null and can be modified by the function, so it can be used to hold the return value if desired, though that is not required. `$0` can also be given an alias. For example, this function works on any data type that has a `+` operator:

```sql
CREATE FUNCTION add_three_values(v1 anyelement, v2 anyelement, v3 anyelement)
RETURNS anyelement AS $$
DECLARE
  result ALIAS FOR $0;
BEGIN
  result := v1 + v2 + v3;
  RETURN result;
END;
$$ LANGUAGE plpgsql;
```
Chapter 41. PL/pgSQL - SQL Procedural Language

The same effect can be obtained by declaring one or more output parameters as polymorphic types. In this case the special $0 parameter is not used; the output parameters themselves serve the same purpose. For example:

```sql
CREATE FUNCTION add_three_values(v1 anyelement, v2 anyelement, v3 anyelement,
                                   OUT sum anyelement)
AS $$
BEGIN
    sum := v1 + v2 + v3;
END;
$$ LANGUAGE plpgsql;
```

### 41.3.2. ALIAS

```sql
newname ALIAS FOR oldname;
```

The ALIAS syntax is more general than is suggested in the previous section: you can declare an alias for any variable, not just function parameters. The main practical use for this is to assign a different name for variables with predetermined names, such as `NEW` or `OLD` within a trigger procedure.

Examples:

```sql
DECLARE
    prior ALIAS FOR old;
    updated ALIAS FOR new;
```

Since ALIAS creates two different ways to name the same object, unrestricted use can be confusing. It’s best to use it only for the purpose of overriding predetermined names.

### 41.3.3. Copying Types

```sql
variable%TYPE
```

%TYPE provides the data type of a variable or table column. You can use this to declare variables that will hold database values. For example, let’s say you have a column named `user_id` in your `users` table. To declare a variable with the same data type as `users.user_id` you write:

```sql
user_id users.user_id%TYPE;
```

By using %TYPE you don’t need to know the data type of the structure you are referencing, and most importantly, if the data type of the referenced item changes in the future (for instance: you change the type of `user_id` from integer to real), you might not need to change your function definition. %TYPE is particularly valuable in polymorphic functions, since the data types needed for internal variables can change from one call to the next. Appropriate variables can be created by applying %TYPE to the function’s arguments or result placeholders.
41.3.4. Row Types

name table_name%ROWTYPE;
name composite_type_name;

A variable of a composite type is called a row variable (or row-type variable). Such a variable can hold a whole row of a `SELECT` or `FOR` query result, so long as that query’s column set matches the declared type of the variable. The individual fields of the row value are accessed using the usual dot notation, for example `rowvar.field`.

A row variable can be declared to have the same type as the rows of an existing table or view, by using the `table_name%ROWTYPE` notation; or it can be declared by giving a composite type’s name. (Since every table has an associated composite type of the same name, it actually does not matter in PostgreSQL whether you write `%ROWTYPE` or not. But the form with `%ROWTYPE` is more portable.)

Parameters to a function can be composite types (complete table rows). In that case, the corresponding identifier $n$ will be a row variable, and fields can be selected from it, for example `$1.user_id`.

Only the user-defined columns of a table row are accessible in a row-type variable, not the OID or other system columns (because the row could be from a view). The fields of the row type inherit the table’s field size or precision for data types such as `char(n)`.

Here is an example of using composite types. `table1` and `table2` are existing tables having at least the mentioned fields:

```sql
CREATE FUNCTION merge_fields(t_row table1) RETURNS text AS $$
DECLARE
    t2_row table2%ROWTYPE;
BEGIN
    SELECT * INTO t2_row FROM table2 WHERE ... ;
    RETURN t_row.f1 || t2_row.f3 || t_row.f5 || t2_row.f7;
END;
$$ LANGUAGE plpgsql;

SELECT merge_fields(t.*) FROM table1 t WHERE ... ;
```

41.3.5. Record Types

name RECORD;

Record variables are similar to row-type variables, but they have no predefined structure. They take on the actual row structure of the row they are assigned during a `SELECT` or `FOR` command. The substructure of a record variable can change each time it is assigned to. A consequence of this is that until a record variable is first assigned to, it has no substructure, and any attempt to access a field in it will draw a run-time error.

Note that `RECORD` is not a true data type, only a placeholder. One should also realize that when a PL/pgSQL function is declared to return type `record`, this is not quite the same concept as a record variable, even though such a function might use a record variable to hold its result. In both cases the actual row structure is unknown when the function is written, but for a function returning `record` the actual structure is determined when the calling query is parsed, whereas a record variable can change its row structure on-the-fly.
41.3.6. Collation of PL/pgSQL Variables

When a PL/pgSQL function has one or more parameters of collatable data types, a collation is identified for each function call depending on the collations assigned to the actual arguments, as described in Section 23.2. If a collation is successfully identified (i.e., there are no conflicts of implicit collations among the arguments) then all the collatable parameters are treated as having that collation implicitly. This will affect the behavior of collation-sensitive operations within the function. For example, consider

```
CREATE FUNCTION less_than(a text, b text) RETURNS boolean AS $$
BEGIN
  RETURN a < b;
END;
$$ LANGUAGE plpgsql;
```

SELECT less_than(text_field_1, text_field_2) FROM table1;
SELECT less_than(text_field_1, text_field_2 COLLATE "C") FROM table1;

The first use of `less_than` will use the common collation of `text_field_1` and `text_field_2` for the comparison, while the second use will use `C` collation.

Furthermore, the identified collation is also assumed as the collation of any local variables that are of collatable types. Thus this function would not work any differently if it were written as

```
CREATE FUNCTION less_than(a text, b text) RETURNS boolean AS $$
DECLARE
  local_a text := a;
  local_b text := b;
BEGIN
  RETURN local_a < local_b;
END;
$$ LANGUAGE plpgsql;
```

If there are no parameters of collatable data types, or no common collation can be identified for them, then parameters and local variables use the default collation of their data type (which is usually the database’s default collation, but could be different for variables of domain types).

A local variable of a collatable data type can have a different collation associated with it by including the `COLLATE` option in its declaration, for example

```
DECLARE
  local_a text COLLATE "en_US";
```

This option overrides the collation that would otherwise be given to the variable according to the rules above.

Also, of course explicit `COLLATE` clauses can be written inside a function if it is desired to force a particular collation to be used in a particular operation. For example,

```
CREATE FUNCTION less_than_c(a text, b text) RETURNS boolean AS $$
BEGIN
  RETURN a < b COLLATE "C";
END;
$$ LANGUAGE plpgsql;
```

This overrides the collations associated with the table columns, parameters, or local variables used in the expression, just as would happen in a plain SQL command.
41.4. Expressions

All expressions used in PL/pgSQL statements are processed using the server’s main SQL executor. For example, when you write a PL/pgSQL statement like

```sql
IF expression THEN ...
```

PL/pgSQL will evaluate the expression by feeding a query like

```sql
SELECT expression
```

to the main SQL engine. While forming the `SELECT` command, any occurrences of PL/pgSQL variable names are replaced by parameters, as discussed in detail in Section 41.10.1. This allows the query plan for the `SELECT` to be prepared just once and then reused for subsequent evaluations with different values of the variables. Thus, what really happens on first use of an expression is essentially a `PREPARE` command. For example, if we have declared two integer variables `x` and `y`, and we write

```sql
IF x < y THEN ...
```

what happens behind the scenes is equivalent to

```sql
PREPARE statement_name(integer, integer) AS SELECT $1 < $2;
```

and then this prepared statement is executed for each execution of the `IF` statement, with the current values of the PL/pgSQL variables supplied as parameter values. Normally these details are not important to a PL/pgSQL user, but they are useful to know when trying to diagnose a problem. More information appears in Section 41.10.2.

41.5. Basic Statements

In this section and the following ones, we describe all the statement types that are explicitly understood by PL/pgSQL. Anything not recognized as one of these statement types is presumed to be an SQL command and is sent to the main database engine to execute, as described in Section 41.5.2 and Section 41.5.3.

41.5.1. Assignment

An assignment of a value to a PL/pgSQL variable is written as:

```sql
variable { := | = } expression;
```

As explained previously, the expression in such a statement is evaluated by means of an SQL `SELECT` command sent to the main database engine. The expression must yield a single value (possibly a row value, if the variable is a row or record variable). The target variable can be a simple variable (optionally qualified with a block name), a field of a row or record variable, or an element of an array that is a simple variable or field. Equal (=) can be used instead of PL/SQL-compliant :=.

If the expression’s result data type doesn’t match the variable’s data type, the value will be coerced as though by an assignment cast (see Section 10.4). If no assignment cast is known for the pair of data types involved, the PL/pgSQL interpreter will attempt to convert the result value textually, that is by applying the result type’s output function followed by the variable type’s input function. Note that this could result in run-time errors generated by the input function, if the string form of the result value is not acceptable to the input function.
Examples:

tax := subtotal * 0.06;
my_record.user_id := 20;

41.5.2. Executing a Command With No Result

For any SQL command that does not return rows, for example INSERT without a RETURNING clause, you can execute the command within a PL/pgSQL function just by writing the command.

Any PL/pgSQL variable name appearing in the command text is treated as a parameter, and then the current value of the variable is provided as the parameter value at run time. This is exactly like the processing described earlier for expressions; for details see Section 41.10.1.

When executing a SQL command in this way, PL/pgSQL may cache and re-use the execution plan for the command, as discussed in Section 41.10.2.

Sometimes it is useful to evaluate an expression or SELECT query but discard the result, for example when calling a function that has side-effects but no useful result value. To do this in PL/pgSQL, use the PERFORM statement:

PERFORM query;

This executes query and discards the result. Write the query the same way you would write an SQL SELECT command, but replace the initial keyword SELECT with PERFORM. For WITH queries, use PERFORM and then place the query in parentheses. (In this case, the query can only return one row.) PL/pgSQL variables will be substituted into the query just as for commands that return no result, and the plan is cached in the same way. Also, the special variable FOUND is set to true if the query produced at least one row, or false if it produced no rows (see Section 41.5.5).

Note: One might expect that writing SELECT directly would accomplish this result, but at present the only accepted way to do it is PERFORM. A SQL command that can return rows, such as SELECT, will be rejected as an error unless it has an INTO clause as discussed in the next section.

An example:

PERFORM create_mv('cs_session_page_requests_mv', my_query);

41.5.3. Executing a Query with a Single-row Result

The result of a SQL command yielding a single row (possibly of multiple columns) can be assigned to a record variable, row-type variable, or list of scalar variables. This is done by writing the base SQL command and adding an INTO clause. For example,

SELECT select_expressions INTO [STRICT] target FROM ...;
INSERT ... RETURNING expressions INTO [STRICT] target;
UPDATE ... RETURNING expressions INTO [STRICT] target;
DELETE ... RETURNING expressions INTO [STRICT] target;
where `target` can be a record variable, a row variable, or a comma-separated list of simple variables and record/row fields. PL/pgSQL variables will be substituted into the rest of the query, and the plan is cached, just as described above for commands that do not return rows. This works for `SELECT, INSERT/UPDATE/DELETE` with `RETURNING`, and utility commands that return row-set results (such as `EXPLAIN`). Except for the `INTO` clause, the SQL command is the same as it would be written outside PL/pgSQL.

**Tip:** Note that this interpretation of `SELECT` with `INTO` is quite different from PostgreSQL’s regular `SELECT INTO` command, wherein the `INTO` target is a newly created table. If you want to create a table from a `SELECT` result inside a PL/pgSQL function, use the syntax `CREATE TABLE ... AS SELECT`.

If a row or a variable list is used as target, the query’s result columns must exactly match the structure of the target as to number and data types, or else a run-time error occurs. When a record variable is the target, it automatically configures itself to the row type of the query result columns.

The `INTO` clause can appear almost anywhere in the SQL command. Customarily it is written either just before or just after the list of `select_expressions` in a `SELECT` command, or at the end of the command for other command types. It is recommended that you follow this convention in case the PL/pgSQL parser becomes stricter in future versions.

If `STRICT` is not specified in the `INTO` clause, then `target` will be set to the first row returned by the query, or to nulls if the query returned no rows. (Note that “the first row” is not well-defined unless you’ve used `ORDER BY`.) Any result rows after the first row are discarded. You can check the special `FOUND` variable (see Section 41.5.5) to determine whether a row was returned:

```sql
SELECT * INTO myrec FROM emp WHERE empname = myname;
IF NOT FOUND THEN
  RAISE EXCEPTION 'employee % not found', myname;
END IF;
```

If the `STRICT` option is specified, the query must return exactly one row or a run-time error will be reported, either `NO_DATA_FOUND` (no rows) or `TOO_MANY_ROWS` (more than one row). You can use an exception block if you wish to catch the error, for example:

```sql
BEGIN
  SELECT * INTO STRICT myrec FROM emp WHERE empname = myname;
EXCEPTION
    WHEN NO_DATA_FOUND THEN
      RAISE EXCEPTION 'employee % not found', myname;
    WHEN TOO_MANY_ROWS THEN
      RAISE EXCEPTION 'employee % not unique', myname;
END;
```

Successful execution of a command with `STRICT` always sets `FOUND` to `true`.

For `INSERT/UPDATE/DELETE` with `RETURNING`, PL/pgSQL reports an error for more than one returned row, even when `STRICT` is not specified. This is because there is no option such as `ORDER BY` with which to determine which affected row should be returned.

If `print_strict_params` is enabled for the function, then when an error is thrown because the requirements of `STRICT` are not met, the `DETAIL` part of the error message will include information about the parameters passed to the query. You can change the `print_strict_params` setting for all functions by setting `plpgsql.print_strict_params`, though only subsequent function compila-
Chapter 41. PL/pgSQL - SQL Procedural Language

Functions will be affected. You can also enable it on a per-function basis by using a compiler option, for example:

```sql
CREATE FUNCTION get_userid(username text) RETURNS int
AS $$
# print_strict_params on
DECLARE
userid int;
BEGIN
    SELECT users.userid INTO STRICT userid
    FROM users WHERE users.username = get_userid.username;
    RETURN userid;
END $$ LANGUAGE plpgsql;
```

On failure, this function might produce an error message such as

```
ERROR: query returned no rows
DETAIL: parameters: $1 = 'nosuchuser'
CONTEXT: PL/pgSQL function get_userid(text) line 6 at SQL statement
```

**Note:** The **STRICT** option matches the behavior of Oracle PL/SQL's **SELECT INTO** and related statements.

To handle cases where you need to process multiple result rows from a SQL query, see Section 41.6.4.

### 41.5.4. Executing Dynamic Commands

Oftentimes you will want to generate dynamic commands inside your PL/pgSQL functions, that is, commands that will involve different tables or different data types each time they are executed. PL/pgSQL's normal attempts to cache plans for commands (as discussed in Section 41.10.2) will not work in such scenarios. To handle this sort of problem, the **EXECUTE** statement is provided:

```
EXECUTE command-string [ INTO [STRICT] target ] [ USING expression [, ... ] ];
```

where **command-string** is an expression yielding a string (of type **text**) containing the command to be executed. The optional **target** is a record variable, a row variable, or a comma-separated list of simple variables and record/row fields, into which the results of the command will be stored. The optional **USING** expressions supply values to be inserted into the command.

No substitution of PL/pgSQL variables is done on the computed command string. Any required variable values must be inserted in the command string as it is constructed; or you can use parameters as described below.

Also, there is no plan caching for commands executed via **EXECUTE**. Instead, the command is always planned each time the statement is run. Thus the command string can be dynamically created within the function to perform actions on different tables and columns.

The **INTO** clause specifies where the results of a SQL command returning rows should be assigned. If a row or variable list is provided, it must exactly match the structure of the query’s results (when a record variable is used, it will configure itself to match the result structure automatically). If multiple
rows are returned, only the first will be assigned to the INTO variable. If no rows are returned, NULL is assigned to the INTO variable(s). If no INTO clause is specified, the query results are discarded.

If the STRICT option is given, an error is reported unless the query produces exactly one row.

The command string can use parameter values, which are referenced in the command as $1, $2, etc. These symbols refer to values supplied in the USING clause. This method is often preferable to inserting data values into the command string as text: it avoids run-time overhead of converting the values to text and back, and it is much less prone to SQL-injection attacks since there is no need for quoting or escaping. An example is:

```sql
EXECUTE 'SELECT count(*) FROM mytable WHERE inserted_by = $1 AND inserted <= $2'
INTO c
USING checked_user, checked_date;
```

Note that parameter symbols can only be used for data values — if you want to use dynamically determined table or column names, you must insert them into the command string textually. For example, if the preceding query needed to be done against a dynamically selected table, you could do this:

```sql
EXECUTE 'SELECT count(*) FROM ' || quote_ident(tabname) || ' WHERE inserted_by = $1 AND inserted <= $2'
INTO c
USING checked_user, checked_date;
```

A cleaner approach is to use format()'s %I specification for table or column names (strings separated by a newline are concatenated):

```sql
EXECUTE format('SELECT count(*) FROM %I ' 'WHERE inserted_by = $1 AND inserted <= $2', tabname)
INTO c
USING checked_user, checked_date;
```

Another restriction on parameter symbols is that they only work in SELECT, INSERT, UPDATE, and DELETE commands. In other statement types (generically called utility statements), you must insert values textually even if they are just data values.

An EXECUTE with a simple constant command string and some USING parameters, as in the first example above, is functionally equivalent to just writing the command directly in PL/pgSQL and allowing replacement of PL/pgSQL variables to happen automatically. The important difference is that EXECUTE will re-plan the command on each execution, generating a plan that is specific to the current parameter values; whereas PL/pgSQL may otherwise create a generic plan and cache it for re-use. In situations where the best plan depends strongly on the parameter values, it can be helpful to use EXECUTE to positively ensure that a generic plan is not selected.

SELECT INTO is not currently supported within EXECUTE; instead, execute a plain SELECT command and specify INTO as part of the EXECUTE itself.

**Note:** The PL/pgSQL EXECUTE statement is not related to the EXECUTE SQL statement supported by the PostgreSQL server. The server's EXECUTE statement cannot be used directly within PL/pgSQL functions (and is not needed).
Example 41-1. Quoting Values In Dynamic Queries

When working with dynamic commands you will often have to handle escaping of single quotes. The recommended method for quoting fixed text in your function body is dollar quoting. (If you have legacy code that does not use dollar quoting, please refer to the overview in Section 41.11.1, which can save you some effort when translating said code to a more reasonable scheme.)

Dynamic values require careful handling since they might contain quote characters. An example using `format()` (this assumes that you are dollar quoting the function body so quote marks need not be doubled):

```sql
EXECUTE format('UPDATE tbl SET %I = $1 
' 'WHERE key = $2', colname) USING newvalue, keyvalue;
```

It is also possible to call the quoting functions directly:

```sql
EXECUTE 'UPDATE tbl SET ' || quote_ident(colname) || ' = ' || quote_literal(newvalue) || ' WHERE key = ' || quote_literal(keyvalue);
```

This example demonstrates the use of the `quote_ident` and `quote_literal` functions (see Section 9.4). For safety, expressions containing column or table identifiers should be passed through `quote_ident` before insertion in a dynamic query. Expressions containing values that should be literal strings in the constructed command should be passed through `quote_literal`. These functions take the appropriate steps to return the input text enclosed in double or single quotes respectively, with any embedded special characters properly escaped.

Because `quote_literal` is labeled `STRICT`, it will always return null when called with a null argument. In the above example, if `newvalue` or `keyvalue` were null, the entire dynamic query string would become null, leading to an error from `EXECUTE`. You can avoid this problem by using the `quote_nullable` function, which works the same as `quote_literal` except that when called with a null argument it returns the string `NULL`. For example,

```sql
EXECUTE 'UPDATE tbl SET ' || quote_ident(colname) || ' = ' || quote_nullable(newvalue) || ' WHERE key = ' || quote_nullable(keyvalue);
```

If you are dealing with values that might be null, you should usually use `quote_nullable` in place of `quote_literal`.

As always, care must be taken to ensure that null values in a query do not deliver unintended results. For example the `WHERE` clause

```sql
WHERE key = ’ || quote_nullable(keyvalue)
```

will never succeed if `keyvalue` is null, because the result of using the equality operator `=` with a null operand is always null. If you wish null to work like an ordinary key value, you would need to rewrite the above as

```sql
WHERE key IS NOT DISTINCT FROM ’ || quote_nullable(keyvalue)
```

(At present, `IS NOT DISTINCT FROM` is handled much less efficiently than `=`, so don’t do this unless you must. See Section 9.2 for more information on nulls and `IS DISTINCT`.)

Note that dollar quoting is only useful for quoting fixed text. It would be a very bad idea to try to write this example as:

```sql
EXECUTE ‘UPDATE tbl SET ’ || quote_ident(colname)
```
because it would break if the contents of newvalue happened to contain $$$. The same objection would apply to any other dollar-quoting delimiter you might pick. So, to safely quote text that is not known in advance, you must use quote_literal, quote_nullable, or quote_ident, as appropriate.

Dynamic SQL statements can also be safely constructed using the format function (see Section 9.4). For example:

```sql
EXECUTE format('UPDATE tbl SET %I = %L ' 'WHERE key = %L', colname, newvalue, keyvalue);
```

%I is equivalent to quote_ident, and %L is equivalent to quote_nullable. The format function can be used in conjunction with the using clause:

```sql
EXECUTE format('UPDATE tbl SET %I = $1 WHERE key = $2', colname) USING newvalue, keyvalue;
```

This form is better because the variables are handled in their native data type format, rather than unconditionally converting them to text and quoting them via %L. It is also more efficient.

A much larger example of a dynamic command and EXECUTE can be seen in Example 41-9, which builds and executes a CREATE FUNCTION command to define a new function.

### 41.5.5. Obtaining the Result Status

There are several ways to determine the effect of a command. The first method is to use the GET DIAGNOSTICS command, which has the form:

```
GET [ CURRENT ] DIAGNOSTICS variable { = | := } item [ , ... ];
```

This command allows retrieval of system status indicators. CURRENT is a noise word (but see also GET STACKED DIAGNOSTICS in Section 41.6.6.1). Each item is a key word identifying a status value to be assigned to the specified variable (which should be of the right data type to receive it). The currently available status items are shown in Table 41-1. Colon-equal (:=) can be used instead of the SQL-standard = token. An example:

```
GET DIAGNOSTICS integer_var = ROW_COUNT;
```

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ROW_COUNT</td>
<td>bigint</td>
<td>the number of rows processed by the most recent SQL command</td>
</tr>
<tr>
<td>RESULT_OID</td>
<td>oid</td>
<td>the OID of the last row inserted by the most recent SQL command (only useful after an INSERT command into a table having OIDs)</td>
</tr>
</tbody>
</table>
The second method to determine the effects of a command is to check the special variable named `FOUND`, which is of type boolean. `FOUND` starts out false within each PL/pgSQL function call. It is set by each of the following types of statements:

- A `SELECT INTO` statement sets `FOUND` true if a row is assigned, false if no row is returned.
- A `PERFORM` statement sets `FOUND` true if it produces (and discards) one or more rows, false if no row is produced.
- `UPDATE`, `INSERT`, and `DELETE` statements set `FOUND` true if at least one row is affected, false if no row is affected.
- A `FETCH` statement sets `FOUND` true if it returns a row, false if no row is returned.
- A `MOVE` statement sets `FOUND` true if it successfully repositions the cursor, false otherwise.
- A `FOR` or `FOREACH` statement sets `FOUND` true if it iterates one or more times, else false. `FOUND` is set this way when the loop exits; inside the execution of the loop, `FOUND` is not modified by the loop statement, although it might be changed by the execution of other statements within the loop body.
- `RETURN QUERY` and `RETURN QUERY EXECUTE` statements set `FOUND` true if the query returns at least one row, false if no row is returned.

Other PL/pgSQL statements do not change the state of `FOUND`. Note in particular that `EXECUTE` changes the output of `GET DIAGNOSTICS`, but does not change `FOUND`.

`FOUND` is a local variable within each PL/pgSQL function; any changes to it affect only the current function.

### 41.5.6. Doing Nothing At All

Sometimes a placeholder statement that does nothing is useful. For example, it can indicate that one arm of an `if/then/else` chain is deliberately empty. For this purpose, use the `NULL` statement:

```
NULL;
```

For example, the following two fragments of code are equivalent:

```
BEGIN
    y := x / 0;
EXCEPTION
    WHEN division_by_zero THEN
        NULL; -- ignore the error
END;

BEGIN
    y := x / 0;
EXCEPTION
    WHEN division_by_zero THEN -- ignore the error
```
Which is preferable is a matter of taste.

Note: In Oracle's PL/SQL, empty statement lists are not allowed, and so NULL statements are required for situations such as this. PL/pgSQL allows you to just write nothing, instead.

41.6. Control Structures

Control structures are probably the most useful (and important) part of PL/pgSQL. With PL/pgSQL’s control structures, you can manipulate PostgreSQL data in a very flexible and powerful way.

41.6.1. Returning From a Function

There are two commands available that allow you to return data from a function: RETURN and RETURN NEXT.

41.6.1.1. RETURN

RETURN expression;

RETURN with an expression terminates the function and returns the value of expression to the caller. This form is used for PL/pgSQL functions that do not return a set.

In a function that returns a scalar type, the expression’s result will automatically be cast into the function’s return type as described for assignments. But to return a composite (row) value, you must write an expression delivering exactly the requested column set. This may require use of explicit casting.

If you declared the function with output parameters, write just RETURN with no expression. The current values of the output parameter variables will be returned.

If you declared the function to return void, a RETURN statement can be used to exit the function early; but do not write an expression following RETURN.

The return value of a function cannot be left undefined. If control reaches the end of the top-level block of the function without hitting a RETURN statement, a run-time error will occur. This restriction does not apply to functions with output parameters and functions returning void, however. In those cases a RETURN statement is automatically executed if the top-level block finishes.

Some examples:

-- functions returning a scalar type
RETURN 1 + 2;
RETURN scalar_var;

-- functions returning a composite type
RETURN composite_type_var;
RETURN (1, 2, 'three'::text); -- must cast columns to correct types
41.6.1.2. **RETURN NEXT and RETURN QUERY**

```
RETURN NEXT expression;
RETURN QUERY query;
RETURN QUERY EXECUTE command-string [ USING expression [, ... ] ];
```

When a PL/pgSQL function is declared to return `SETOF sometype`, the procedure to follow is slightly different. In that case, the individual items to return are specified by a sequence of `RETURN NEXT` or `RETURN QUERY` commands, and then a final `RETURN` command with no argument is used to indicate that the function has finished executing. `RETURN NEXT` can be used with both scalar and composite data types; with a composite result type, an entire “table” of results will be returned. `RETURN QUERY` appends the results of executing a query to the function’s result set. `RETURN NEXT` and `RETURN QUERY` can be freely intermixed in a single set-returning function, in which case their results will be concatenated.

`RETURN NEXT` and `RETURN QUERY` do not actually return from the function — they simply append zero or more rows to the function’s result set. Execution then continues with the next statement in the PL/pgSQL function. As successive `RETURN NEXT` or `RETURN QUERY` commands are executed, the result set is built up. A final `RETURN`, which should have no argument, causes control to exit the function (or you can just let control reach the end of the function).

`RETURN QUERY` has a variant `RETURN QUERY EXECUTE`, which specifies the query to be executed dynamically. Parameter expressions can be inserted into the computed query string via `USING`, in just the same way as in the `EXECUTE` command.

If you declared the function with output parameters, write just `RETURN NEXT` with no expression. On each execution, the current values of the output parameter variable(s) will be saved for eventual return as a row of the result. Note that you must declare the function as returning `SETOF record` when there are multiple output parameters, or `SETOF sometype` when there is just one output parameter of type `sometype`, in order to create a set-returning function with output parameters.

Here is an example of a function using `RETURN NEXT`:

```
CREATE TABLE foo (fooid INT, foosubid INT, fooname TEXT);
INSERT INTO foo VALUES (1, 2, 'three');
INSERT INTO foo VALUES (4, 5, 'six');

CREATE OR REPLACE FUNCTION get_all_foo() RETURNS SETOF foo AS $BODY$
DECLARE
    r foo%rowtype;
BEGIN
    FOR r IN
        SELECT * FROM foo WHERE fooid > 0
    LOOP
        -- can do some processing here
        RETURN NEXT r; -- return current row of SELECT
    END LOOP;
    RETURN;
END
$BODY$
LANGUAGE plpgsql;

SELECT * FROM get_all_foo();
```

Here is an example of a function using `RETURN QUERY`:

```
CREATE TABLE foo (fooid INT, foosubid INT, fooname TEXT);
INSERT INTO foo VALUES (1, 2, 'three');
INSERT INTO foo VALUES (4, 5, 'six');

CREATE OR REPLACE FUNCTION get_all_foo() RETURNS SETOF foo AS $BODY$
DECLARE
    r foo%rowtype;
BEGIN
    FOR r IN
        SELECT * FROM foo WHERE fooid > 0
    LOOP
        -- can do some processing here
        RETURN QUERY EXECUTE 'SELECT * FROM foo WHERE fooid = $1';
    END LOOP;
    RETURN;
END
$BODY$
LANGUAGE plpgsql;

SELECT * FROM get_all_foo();
```
CREATE FUNCTION get_available_flightid(date) RETURNS SETOF integer AS
$$BODY$
BEGIN
RETURN QUERY SELECT flightid
FROM flight
WHERE flightdate >= $1
AND flightdate < ($1 + 1);

-- Since execution is not finished, we can check whether rows were returned
-- and raise exception if not.
IF NOT FOUND THEN
RAISE EXCEPTION 'No flight at %.', $1;
END IF;

RETURN;
END
$$BODY$
LANGUAGE plpgsql;

-- Returns available flights or raises exception if there are no
-- available flights.
SELECT * FROM get_available_flightid(CURRENT_DATE);

Note: The current implementation of RETURN NEXT and RETURN QUERY stores the entire result set
before returning from the function, as discussed above. That means that if a PL/pgSQL function
produces a very large result set, performance might be poor: data will be written to disk to avoid
memory exhaustion, but the function itself will not return until the entire result set has been gener-
ated. A future version of PL/pgSQL might allow users to define set-returning functions that do not
have this limitation. Currently, the point at which data begins being written to disk is controlled by
the work_mem configuration variable. Administrators who have sufficient memory to store larger
result sets in memory should consider increasing this parameter.

41.6.2. Conditionals

IF and CASE statements let you execute alternative commands based on certain conditions. PL/pgSQL
has three forms of IF:

- IF ... THEN ... END IF
- IF ... THEN ... ELSE ... END IF
- IF ... THEN ... ELSIF ... THEN ... ELSE ... END IF

and two forms of CASE:

- CASE ... WHEN ... THEN ... ELSE ... END CASE
- CASE WHEN ... THEN ... ELSE ... END CASE
41.6.2.1. IF–THEN

**IF** boolean-expression **THEN**

  statements

**END IF;**

IF–THEN statements are the simplest form of IF. The statements between **THEN** and **END IF** will be executed if the condition is true. Otherwise, they are skipped.

Example:

```sql
IF v_user_id <> 0 THEN
    UPDATE users SET email = v_email WHERE user_id = v_user_id;
**END IF;**
```

41.6.2.2. IF–THEN–ELSE

**IF** boolean-expression **THEN**

  statements

**ELSE**

  statements

**END IF;**

IF–THEN–ELSE statements add to IF–THEN by letting you specify an alternative set of statements that should be executed if the condition is not true. (Note this includes the case where the condition evaluates to NULL.)

Examples:

```sql
IF parentid IS NULL OR parentid = "
THEN
    RETURN fullname;
ELSE
    RETURN hp_true_filename(parentid) || '/' || fullname;
**END IF;**

IF v_count > 0 THEN
    INSERT INTO users_count (count) VALUES (v_count);
    RETURN 't';
ELSE
    RETURN 'f';
**END IF;**
```

41.6.2.3. IF–THEN–ELSIF

**IF** boolean-expression **THEN**

  statements

[ **ELSIF** boolean-expression **THEN**

  statements

  ...]]

...
Sometimes there are more than just two alternatives. IF-THEN-ELSIF provides a convenient method of checking several alternatives in turn. The IF conditions are tested successively until the first one that is true is found. Then the associated statement(s) are executed, after which control passes to the next statement after END IF. (Any subsequent IF conditions are not tested.) If none of the IF conditions is true, then the ELSE block (if any) is executed.

Here is an example:

```sql
IF number = 0 THEN
  result := 'zero';
ELSIF number > 0 THEN
  result := 'positive';
ELSIF number < 0 THEN
  result := 'negative';
ELSE
  -- hmm, the only other possibility is that number is null
  result := 'NULL';
END IF;
```

The key word ELSIF can also be spelled ELSEIF.

An alternative way of accomplishing the same task is to nest IF-THEN-ELSE statements, as in the following example:

```sql
IF demo_row.sex = 'm' THEN
  pretty_sex := 'man';
ELSE
  IF demo_row.sex = 'f' THEN
    pretty_sex := 'woman';
  END IF;
END IF;
```

However, this method requires writing a matching END IF for each IF, so it is much more cumbersome than using ELSIF when there are many alternatives.

### 41.6.2.4. Simple CASE

```sql
CASE search-expression
  WHEN expression [, expression [ ... ]] THEN
    statements
  [ WHEN expression [, expression [ ... ]] THEN
    statements
  ... ]
  [ ELSE
    statements ]
END CASE;
```

The simple form of CASE provides conditional execution based on equality of operands. The search-expression is evaluated (once) and successively compared to each expression in the
WHEN clauses. If a match is found, then the corresponding statements are executed, and then control passes to the next statement after END CASE. (Subsequent WHEN expressions are not evaluated.) If no match is found, the ELSE statements are executed; but if ELSE is not present, then a CASE_NOT_FOUND exception is raised.

Here is a simple example:

```sql
CASE x
  WHEN 1, 2 THEN
    msg := 'one or two';
  ELSE
    msg := 'other value than one or two';
END CASE;
```

### 41.6.2.5. Searched CASE

CASE
  WHEN boolean-expression THEN
    statements
  [ WHEN boolean-expression THEN
    statements
  ... ]
  [ ELSE
    statements
  ]
END CASE;

The searched form of CASE provides conditional execution based on truth of Boolean expressions. Each WHEN clause’s boolean-expression is evaluated in turn, until one is found that yields true. Then the corresponding statements are executed, and then control passes to the next statement after END CASE. (Subsequent WHEN expressions are not evaluated.) If no true result is found, the ELSE statements are executed; but if ELSE is not present, then a CASE_NOT_FOUND exception is raised.

Here is an example:

```sql
CASE
  WHEN x BETWEEN 0 AND 10 THEN
    msg := 'value is between zero and ten';
  WHEN x BETWEEN 11 AND 20 THEN
    msg := 'value is between eleven and twenty';
END CASE;
```

This form of CASE is entirely equivalent to IF-THEN-ELSIF, except for the rule that reaching an omitted ELSE clause results in an error rather than doing nothing.

### 41.6.3. Simple Loops

With the LOOP, EXIT, CONTINUE, WHILE, FOR, and FOREACH statements, you can arrange for your PL/pgSQL function to repeat a series of commands.
41.6.3.1. LOOP

[ <label>> ]
LOOP
  statements
END LOOP [ label ];

LOOP defines an unconditional loop that is repeated indefinitely until terminated by an EXIT or RETURN statement. The optional label can be used by EXIT and CONTINUE statements within nested loops to specify which loop those statements refer to.

41.6.3.2. EXIT

EXIT [ label ] [ WHEN boolean-expression ];

If no label is given, the innermost loop is terminated and the statement following END LOOP is executed next. If label is given, it must be the label of the current or some outer level of nested loop or block. Then the named loop or block is terminated and control continues with the statement after the loop's/block's corresponding END.

If WHEN is specified, the loop exit occurs only if boolean-expression is true. Otherwise, control passes to the statement after EXIT.

EXIT can be used with all types of loops; it is not limited to use with unconditional loops.

When used with a BEGIN block, EXIT passes control to the next statement after the end of the block. Note that a label must be used for this purpose; an unlabeled EXIT is never considered to match a BEGIN block. (This is a change from pre-8.4 releases of PostgreSQL, which would allow an unlabeled EXIT to match a BEGIN block.)

Examples:

LOOP
  -- some computations
  IF count > 0 THEN
    EXIT; -- exit loop
  END IF;
END LOOP;

LOOP
  -- some computations
  EXIT WHEN count > 0; -- same result as previous example
END LOOP;

<<ablock>>
BEGIN
  -- some computations
  IF stocks > 100000 THEN
    EXIT ablock; -- causes exit from the BEGIN block
  END IF;
  -- computations here will be skipped when stocks > 10000
END;
### 41.6.3.3. CONTINUE

CONTINUE [ label ] [ WHEN boolean-expression ];

If no `label` is given, the next iteration of the innermost loop is begun. That is, all statements remaining in the loop body are skipped, and control returns to the loop control expression (if any) to determine whether another loop iteration is needed. If `label` is present, it specifies the label of the loop whose execution will be continued.

If `WHEN` is specified, the next iteration of the loop is begun only if `boolean-expression` is true. Otherwise, control passes to the statement after `CONTINUE`.

`CONTINUE` can be used with all types of loops; it is not limited to use with unconditional loops.

Examples:

```
LOOP
  -- some computations
  EXIT WHEN count > 100;
  CONTINUE WHEN count < 50;
  -- some computations for count IN [50 .. 100]
END LOOP;
```

### 41.6.3.4. WHILE

```
[ <<<label>>> ]
WHILE boolean-expression LOOP
  statements
END LOOP [ label ];
```

The `WHILE` statement repeats a sequence of statements so long as the `boolean-expression` evaluates to true. The expression is checked just before each entry to the loop body.

For example:

```
WHILE amount_owed > 0 AND gift_certificate_balance > 0 LOOP
  -- some computations here
END LOOP;
```

```
WHILE NOT done LOOP
  -- some computations here
END LOOP;
```

### 41.6.3.5. FOR (Integer Variant)

```
[ <<<label>>> ]
FOR name IN [ REVERSE ] expression .. expression [ BY expression ] LOOP
  statements
END LOOP [ label ];
```

This form of `FOR` creates a loop that iterates over a range of integer values. The variable `name` is automatically defined as type `integer` and exists only inside the loop (any existing definition of the
variable name is ignored within the loop). The two expressions giving the lower and upper bound of the range are evaluated once when entering the loop. If the *BY* clause isn’t specified the iteration step is 1, otherwise it’s the value specified in the *BY* clause, which again is evaluated once on loop entry. If *REVERSE* is specified then the step value is subtracted, rather than added, after each iteration.

Some examples of integer *FOR* loops:

```sql
FOR i IN 1..10 LOOP
  -- i will take on the values 1,2,3,4,5,6,7,8,9,10 within the loop
END LOOP;

FOR i IN REVERSE 10..1 LOOP
  -- i will take on the values 10,9,8,7,6,5,4,3,2,1 within the loop
END LOOP;

FOR i IN REVERSE 10..1 BY 2 LOOP
  -- i will take on the values 10,8,6,4 within the loop
END LOOP;
```

If the lower bound is greater than the upper bound (or less than, in the *REVERSE* case), the loop body is not executed at all. No error is raised.

If a *label* is attached to the *FOR* loop then the integer loop variable can be referenced with a qualified name, using that *label*.

### 41.6.4. Looping Through Query Results

Using a different type of *FOR* loop, you can iterate through the results of a query and manipulate that data accordingly. The syntax is:

```sql
[ [<<<label>>>] ]
FOR target IN query LOOP
  statements
END LOOP [ label ];
```

The *target* is a record variable, row variable, or comma-separated list of scalar variables. The *target* is successively assigned each row resulting from the *query* and the loop body is executed for each row. Here is an example:

```sql
CREATE FUNCTION refresh_mviews() RETURNS integer AS $$
DECLARE
  mviews RECORD;
BEGIN
  RAISE NOTICE 'Refreshing all materialized views...';
  FOR mviews IN
    SELECT n.nspname AS mv_schema,
           c.relname AS mv_name,
           pg_catalog.pg_get_userbyid(c.relowner) AS owner
    FROM pg_catalog.pg_class c
    LEFT JOIN pg_catalog.pg_namespace n ON (n.oid = c.relnamespace)
    WHERE c.relkind = 'm'
  ORDER BY 1
  LOOP
```

1141
Chapter 41. PL/pgSQL - SQL Procedural Language

-- Now "mviews" has one record with information about the materialized view
RAISE NOTICE 'Refreshing materialized view %.% (owner: %)...',
quote_ident(mviews.mv_schema),
quote_ident(mviews.mv_name),
quote_ident(mviews.owner);
EXECUTE format('REFRESH MATERIALIZED VIEW %I.%I', mviews.mv_schema, mviews.mv_name);
END LOOP;
RAISE NOTICE 'Done refreshing materialized views.';
RETURN 1;
END;
$$ LANGUAGE plpgsql;

If the loop is terminated by an EXIT statement, the last assigned row value is still accessible after the loop.

The query used in this type of FOR statement can be any SQL command that returns rows to the caller: SELECT is the most common case, but you can also use INSERT, UPDATE, or DELETE with a RETURNING clause. Some utility commands such as EXPLAIN will work too.

PL/pgSQL variables are substituted into the query text, and the query plan is cached for possible re-use, as discussed in detail in Section 41.10.1 and Section 41.10.2.

The FOR-IN-EXECUTE statement is another way to iterate over rows:

[ <<label>> ]
FOR target IN EXECUTE text_expression [ USING expression [ , ... ] ] LOOP
  statements
END LOOP [ label ];

This is like the previous form, except that the source query is specified as a string expression, which is evaluated and replanned on each entry to the FOR loop. This allows the programmer to choose the speed of a preplanned query or the flexibility of a dynamic query, just as with a plain EXECUTE statement. As with EXECUTE, parameter values can be inserted into the dynamic command via USING.

Another way to specify the query whose results should be iterated through is to declare it as a cursor. This is described in Section 41.7.4.

41.6.5. Looping Through Arrays

The FOREACH loop is much like a FOR loop, but instead of iterating through the rows returned by a SQL query, it iterates through the elements of an array value. (In general, FOREACH is meant for looping through components of a composite-valued expression; variants for looping through composites besides arrays may be added in future.) The FOREACH statement to loop over an array is:

[ <<label>> ]
FOREACH target [ SLICE number ] IN ARRAY expression LOOP
  statements
END LOOP [ label ];

Without SLICE, or if SLICE 0 is specified, the loop iterates through individual elements of the array produced by evaluating the expression. The target variable is assigned each element value in
sequence, and the loop body is executed for each element. Here is an example of looping through the elements of an integer array:

```sql
CREATE FUNCTION sum(int[]) RETURNS int8 AS $$
DECLARE
  s int8 := 0;
  x int;
BEGIN
  FOREACH x IN ARRAY $1
  LOOP
    s := s + x;
  END LOOP;
  RETURN s;
END;
$$ LANGUAGE plpgsql;
```

The elements are visited in storage order, regardless of the number of array dimensions. Although the target is usually just a single variable, it can be a list of variables when looping through an array of composite values (records). In that case, for each array element, the variables are assigned from successive columns of the composite value.

With a positive SLICE value, FOREACH iterates through slices of the array rather than single elements. The SLICE value must be an integer constant not larger than the number of dimensions of the array. The target variable must be an array, and it receives successive slices of the array value, where each slice is of the number of dimensions specified by SLICE. Here is an example of iterating through one-dimensional slices:

```sql
CREATE FUNCTION scan_rows(int[]) RETURNS void AS $$
DECLARE
  x int[];
BEGIN
  FOREACH x SLICE 1 IN ARRAY $1
  LOOP
    RAISE NOTICE 'row = %', x;
  END LOOP;
END;
$$ LANGUAGE plpgsql;
```

```sql
SELECT scan_rows(ARRAY[[1,2,3],[4,5,6],[7,8,9],[10,11,12]]);
```

```
NOTICE: row = {1,2,3}
NOTICE: row = {4,5,6}
NOTICE: row = {7,8,9}
NOTICE: row = {10,11,12}
```

### 41.6.6. Trapping Errors

By default, any error occurring in a PL/pgSQL function aborts execution of the function, and indeed of the surrounding transaction as well. You can trap errors and recover from them by using a BEGIN block with an EXCEPTION clause. The syntax is an extension of the normal syntax for a BEGIN block:

```sql
[ [<<label>>] ]
[ DECLARE
  declarations ]
```
BEGIN
    statements
EXCEPTION
    WHEN condition [ OR condition ... ] THEN
        handler_statements
        [ WHEN condition [ OR condition ... ] THEN
            handler_statements
        ... ]
END;

If no error occurs, this form of block simply executes all the statements, and then control passes to the next statement after END. But if an error occurs within the statements, further processing of the statements is abandoned, and control passes to the EXCEPTION list. The list is searched for the first condition matching the error that occurred. If a match is found, the corresponding handler_statements are executed, and then control passes to the next statement after END. If no match is found, the error propagates out as though the EXCEPTION clause were not there at all: the error can be caught by an enclosing block with EXCEPTION, or if there is none it aborts processing of the function.

The condition names can be any of those shown in Appendix A. A category name matches any error within its category. The special condition name OTHERS matches every error type except QUERY_CANCELED and ASSERT_FAILURE. (It is possible, but often unwise, to trap those two error types by name.) Condition names are not case-sensitive. Also, an error condition can be specified by SQLSTATE code; for example these are equivalent:

WHEN division_by_zero THEN ...
WHEN SQLSTATE '22012' THEN ...

If a new error occurs within the selected handler_statements, it cannot be caught by this EXCEPTION clause, but is propagated out. A surrounding EXCEPTION clause could catch it.

When an error is caught by an EXCEPTION clause, the local variables of the PL/pgSQL function remain as they were when the error occurred, but all changes to persistent database state within the block are rolled back. As an example, consider this fragment:

INSERT INTO mytab(firstname, lastname) VALUES('Tom', 'Jones');
BEGIN
    UPDATE mytab SET firstname = 'Joe' WHERE lastname = 'Jones';
    x := x + 1;
    y := x / 0;
EXCEPTION
    WHEN division_by_zero THEN
        RAISE NOTICE 'caught division_by_zero';
        RETURN x;
END;

When control reaches the assignment to y, it will fail with a division_by_zero error. This will be caught by the EXCEPTION clause. The value returned in the RETURN statement will be the incremented value of x, but the effects of the UPDATE command will have been rolled back. The INSERT command preceding the block is not rolled back, however, so the end result is that the database contains Tom Jones not Joe Jones.
Tip: A block containing an EXCEPTION clause is significantly more expensive to enter and exit than a block without one. Therefore, don’t use EXCEPTION without need.

Example 41-2. Exceptions with UPDATE/INSERT

This example uses exception handling to perform either UPDATE or INSERT, as appropriate. It is recommended that applications use INSERT with ON CONFLICT DO UPDATE rather than actually using this pattern. This example serves primarily to illustrate use of PL/pgSQL control flow structures:

```
CREATE TABLE db (a INT PRIMARY KEY, b TEXT);

CREATE FUNCTION merge_db(key INT, data TEXT) RETURNS VOID AS
$$
BEGIN
  LOOP
    -- first try to update the key
    UPDATE db SET b = data WHERE a = key;
    IF found THEN
      RETURN;
    END IF;
    -- not there, so try to insert the key
    -- if someone else inserts the same key concurrently,
    -- we could get a unique-key failure
    BEGIN
      INSERT INTO db(a,b) VALUES (key, data);
      RETURN;
      EXCEPTION WHEN unique_violation THEN
      -- Do nothing, and loop to try the UPDATE again.
    END;
  END LOOP;
END;
$$
LANGUAGE plpgsql;

SELECT merge_db(1, 'david');
SELECT merge_db(1, 'dennis');
```

This coding assumes the unique_violation error is caused by the INSERT, and not by, say, an INSERT in a trigger function on the table. It might also misbehave if there is more than one unique index on the table, since it will retry the operation regardless of which index caused the error. More safety could be had by using the features discussed next to check that the trapped error was the one expected.

### 41.6.6.1. Obtaining Information About an Error

Exception handlers frequently need to identify the specific error that occurred. There are two ways to get information about the current exception in PL/pgSQL: special variables and the GET STACKED DIAGNOSTICS command.

Within an exception handler, the special variable SQLSTATE contains the error code that corresponds to the exception that was raised (refer to Table A-1 for a list of possible error codes). The special variable SQLERRM contains the error message associated with the exception. These variables are undefined outside exception handlers.
Within an exception handler, one may also retrieve information about the current exception by using the \texttt{GET STACKED DIAGNOSTICS} command, which has the form:

\texttt{GET STACKED DIAGNOSTICS variable \{ = | := \} item [ , ... ];}

Each \texttt{item} is a key word identifying a status value to be assigned to the specified \texttt{variable} (which should be of the right data type to receive it). The currently available status items are shown in Table 41-2.

\begin{table}[h]
\centering
\begin{tabular}{|l|l|l|}
\hline
\textbf{Name} & \textbf{Type} & \textbf{Description} \\
\hline
\texttt{RETURNED\_SQLSTATE} & text & the SQLSTATE error code of the exception \\
\hline
\texttt{COLUMN\_NAME} & text & the name of the column related to exception \\
\hline
\texttt{CONSTRAINT\_NAME} & text & the name of the constraint related to exception \\
\hline
\texttt{PG\_DATATYPE\_NAME} & text & the name of the data type related to exception \\
\hline
\texttt{MESSAGE\_TEXT} & text & the text of the exception’s primary message \\
\hline
\texttt{TABLE\_NAME} & text & the name of the table related to exception \\
\hline
\texttt{SCHEMA\_NAME} & text & the name of the schema related to exception \\
\hline
\texttt{PG\_EXCEPTION\_DETAIL} & text & the text of the exception’s detail message, if any \\
\hline
\texttt{PG\_EXCEPTION\_HINT} & text & the text of the exception’s hint message, if any \\
\hline
\texttt{PG\_EXCEPTION\_CONTEXT} & text & line(s) of text describing the call stack at the time of the exception (see Section 41.6.7) \\
\hline
\end{tabular}
\end{table}

If the exception did not set a value for an item, an empty string will be returned.

Here is an example:

\begin{verbatim}
DECLARE
text_var1 text;
text_var2 text;
text_var3 text;
BEGIN
  -- some processing which might cause an exception 
  ...
EXCEPTION WHEN OTHERS THEN
  GET STACKED DIAGNOSTICS text_var1 = MESSAGE_TEXT,
                  text_var2 = PG\_EXCEPTION\_DETAIL,
                  text_var3 = PG\_EXCEPTION\_HINT;
END;
\end{verbatim}
41.6.7. Obtaining Execution Location Information

The GET DIAGNOSTICS command, previously described in Section 41.5.5, retrieves information about current execution state (whereas the GET STACKED DIAGNOSTICS command discussed above reports information about the execution state as of a previous error). Its PG_CONTEXT status item is useful for identifying the current execution location. PG_CONTEXT returns a text string with line(s) of text describing the call stack. The first line refers to the current function and currently executing GET DIAGNOSTICS command. The second and any subsequent lines refer to calling functions further up the call stack. For example:

```sql
CREATE OR REPLACE FUNCTION outer_func() RETURNS integer AS $$
BEGIN
  RETURN inner_func();
END;
$$ LANGUAGE plpgsql;

CREATE OR REPLACE FUNCTION inner_func() RETURNS integer AS $$
DECLARE
  stack text;
BEGIN
  GET DIAGNOSTICS stack = PG_CONTEXT;
  RAISE NOTICE E'--- Call Stack ---
%', stack;
  RETURN 1;
END;
$$ LANGUAGE plpgsql;

SELECT outer_func();
```

GET STACKED DIAGNOSTICS ... PG_EXCEPTION_CONTEXT returns the same sort of stack trace, but describing the location at which an error was detected, rather than the current location.

41.7. Cursors

Rather than executing a whole query at once, it is possible to set up a cursor that encapsulates the query, and then read the query result a few rows at a time. One reason for doing this is to avoid memory overrun when the result contains a large number of rows. (However, PL/pgSQL users do not normally need to worry about that, since FOR loops automatically use a cursor internally to avoid memory problems.) A more interesting usage is to return a reference to a cursor that a function has created, allowing the caller to read the rows. This provides an efficient way to return large row sets from functions.
41.7.1. Declaring Cursor Variables

All access to cursors in PL/pgSQL goes through cursor variables, which are always of the special data type `refcursor`. One way to create a cursor variable is just to declare it as a variable of type `refcursor`. Another way is to use the cursor declaration syntax, which in general is:

```
name [ [ NO ] SCROLL ] CURSOR [ ( arguments ) ] FOR query;
```

(`FOR` can be replaced by `IS` for Oracle compatibility.) If `SCROLL` is specified, the cursor will be capable of scrolling backward; if `NO SCROLL` is specified, backward fetches will be rejected; if neither specification appears, it is query-dependent whether backward fetches will be allowed. `arguments`, if specified, is a comma-separated list of pairs `name datatype` that define names to be replaced by parameter values in the given query. The actual values to substitute for these names will be specified later, when the cursor is opened.

Some examples:

```
DECLARE
curs1 refcursor;
curs2 CURSOR FOR SELECT * FROM tenk1;
curs3 CURSOR (key integer) FOR SELECT * FROM tenk1 WHERE unique1 = key;
```

All three of these variables have the data type `refcursor`, but the first can be used with any query, while the second has a fully specified query already bound to it, and the last has a parameterized query bound to it. (`key` will be replaced by an integer parameter value when the cursor is opened.) The variable `curs1` is said to be `unbound` since it is not bound to any particular query.

41.7.2. Opening Cursors

Before a cursor can be used to retrieve rows, it must be `opened`. (This is the equivalent action to the SQL command `DECLARE CURSOR`.) PL/pgSQL has three forms of the `OPEN` statement, two of which use unbound cursor variables while the third uses a bound cursor variable.

```
Note: Bound cursor variables can also be used without explicitly opening the cursor, via the `FOR` statement described in Section 41.7.4.
```

41.7.2.1. `OPEN FOR query`

```
OPEN unbound_cursorvar [ [ NO ] SCROLL ] FOR query;
```

The cursor variable is opened and given the specified query to execute. The cursor cannot be open already, and it must have been declared as an unbound cursor variable (that is, as a simple `refcursor` variable). The query must be a `SELECT`, or something else that returns rows (such as `EXPLAIN`). The query is treated in the same way as other SQL commands in PL/pgSQL: PL/pgSQL variable names are substituted, and the query plan is cached for possible reuse. When a PL/pgSQL variable is substituted into the cursor query, the value that is substituted is the one it has at the time of the `OPEN`; subsequent changes to the variable will not affect the cursor’s behavior. The `SCROLL` and `NO SCROLL` options have the same meanings as for a bound cursor.

An example:

```
OPEN curs1 FOR SELECT * FROM foo WHERE key = mykey;
```
### 41.7.2.2. OPEN FOR EXECUTE

```sql
OPEN unbound_cursorvar [ [ NO ] SCROLL ] FOR EXECUTE query_string
[ USING expression [, ... ] ];
```

The cursor variable is opened and given the specified query to execute. The cursor cannot be open already, and it must have been declared as an unbound cursor variable (that is, as a simple refcursor variable). The query is specified as a string expression, in the same way as in the EXECUTE command. As usual, this gives flexibility so the query plan can vary from one run to the next (see Section 41.10.2), and it also means that variable substitution is not done on the command string. As with EXECUTE, parameter values can be inserted into the dynamic command via format() and USING. The SCROLL and NO SCROLL options have the same meanings as for a bound cursor.

An example:

```sql
OPEN curs1 FOR EXECUTE format('SELECT * FROM %I WHERE col1 = $1',tabname) USING keyvalue;
```

In this example, the table name is inserted into the query via format(). The comparison value for col1 is inserted via a USING parameter, so it needs no quoting.

### 41.7.2.3. Opening a Bound Cursor

```sql
OPEN bound_cursorvar [ ( [ argument_name := ] argument_value [, ... ] ) ];
```

This form of OPEN is used to open a cursor variable whose query was bound to it when it was declared. The cursor cannot be open already. A list of actual argument value expressions must appear if and only if the cursor was declared to take arguments. These values will be substituted in the query.

The query plan for a bound cursor is always considered cacheable; there is no equivalent of EXECUTE in this case. Notice that SCROLL and NO SCROLL cannot be specified in OPEN, as the cursor’s scrolling behavior was already determined.

Argument values can be passed using either positional or named notation. In positional notation, all arguments are specified in order. In named notation, each argument’s name is specified using := to separate it from the argument expression. Similar to calling functions, described in Section 4.3, it is also allowed to mix positional and named notation.

Examples (these use the cursor declaration examples above):

```sql
OPEN curs2;
OPEN curs3(42);
OPEN curs3(key := 42);
```

Because variable substitution is done on a bound cursor’s query, there are really two ways to pass values into the cursor: either with an explicit argument to OPEN, or implicitly by referencing a PL/pgSQL variable in the query. However, only variables declared before the bound cursor was declared will be substituted into it. In either case the value to be passed is determined at the time of the OPEN. For example, another way to get the same effect as the curs3 example above is

```sql
DECLARE
    key integer;
```
Chapter 41. PL/pgSQL - SQL Procedural Language

```sql
  curs4 CURSOR FOR SELECT * FROM tenk1 WHERE unique1 = key;
BEGIN
  key := 42;
  OPEN curs4;
```

### 41.7.3. Using Cursors

Once a cursor has been opened, it can be manipulated with the statements described here. These manipulations need not occur in the same function that opened the cursor to begin with. You can return a refcursor value out of a function and let the caller operate on the cursor. (Internally, a refcursor value is simply the string name of a so-called portal containing the active query for the cursor. This name can be passed around, assigned to other refcursor variables, and so on, without disturbing the portal.)

All portals are implicitly closed at transaction end. Therefore a refcursor value is usable to reference an open cursor only until the end of the transaction.

#### 41.7.3.1. FETCH

```sql
  FETCH [ direction { FROM | IN } ] cursor INTO target;
```

FETCH retrieves the next row from the cursor into a target, which might be a row variable, a record variable, or a comma-separated list of simple variables, just like SELECT INTO. If there is no next row, the target is set to NULL(s). As with SELECT INTO, the special variable FOUND can be checked to see whether a row was obtained or not.

The `direction` clause can be any of the variants allowed in the SQL FETCH command except the ones that can fetch more than one row; namely, it can be NEXT, PRIOR, FIRST, LAST, ABSOLUTE COUNT, RELATIVE COUNT, FORWARD, or BACKWARD. Omitting `direction` is the same as specifying NEXT. In the forms using a `count`, the `count` can be any integer-valued expression (unlike the SQL FETCH command, which only allows an integer constant). `direction` values that require moving backward are likely to fail unless the cursor was declared or opened with the SCROLL option.

`cursor` must be the name of a refcursor variable that references an open cursor portal.

Examples:

```sql
  FETCH curs1 INTO rowvar;
  FETCH curs2 INTO foo, bar, baz;
  FETCH LAST FROM curs3 INTO x, y;
  FETCH RELATIVE -2 FROM curs4 INTO x;
```

#### 41.7.3.2. MOVE

```sql
  MOVE [ direction { FROM | IN } ] cursor;
```

MOVE repositions a cursor without retrieving any data. MOVE works exactly like the FETCH command, except it only repositions the cursor and does not return the row moved to. As with SELECT INTO, the special variable FOUND can be checked to see whether there was a next row to move to.

1150
Examples:

MOVE curs1;
MOVE LAST FROM curs3;
MOVE RELATIVE -2 FROM curs4;
MOVE FORWARD 2 FROM curs4;

### 41.7.3.3. UPDATE/DELETE WHERE CURRENT OF

**UPDATE** `table` SET ... **WHERE CURRENT OF** `cursor`;
**DELETE** FROM `table` **WHERE CURRENT OF** `cursor`;

When a cursor is positioned on a table row, that row can be updated or deleted using the cursor to identify the row. There are restrictions on what the cursor’s query can be (in particular, no grouping) and it’s best to use **FOR UPDATE** in the cursor. For more information see the DECLARE reference page.

An example:

```
UPDATE foo SET dataval = myval WHERE CURRENT OF curs1;
```

### 41.7.3.4. CLOSE

**CLOSE** `cursor`;

**CLOSE** closes the portal underlying an open cursor. This can be used to release resources earlier than end of transaction, or to free up the cursor variable to be opened again.

An example:

```
CLOSE curs1;
```

### 41.7.3.5. Returning Cursors

PL/pgSQL functions can return cursors to the caller. This is useful to return multiple rows or columns, especially with very large result sets. To do this, the function opens the cursor and returns the cursor name to the caller (or simply opens the cursor using a portal name specified by or otherwise known to the caller). The caller can then fetch rows from the cursor. The cursor can be closed by the caller, or it will be closed automatically when the transaction closes.

The portal name used for a cursor can be specified by the programmer or automatically generated. To specify a portal name, simply assign a string to the `refcursor` variable before opening it. The string value of the `refcursor` variable will be used by **OPEN** as the name of the underlying portal. However, if the `refcursor` variable is null, **OPEN** automatically generates a name that does not conflict with any existing portal, and assigns it to the `refcursor` variable.

**Note:** A bound cursor variable is initialized to the string value representing its name, so that the portal name is the same as the cursor variable name, unless the programmer overrides it by
assignment before opening the cursor. But an unbound cursor variable defaults to the null value initially, so it will receive an automatically-generated unique name, unless overridden.

The following example shows one way a cursor name can be supplied by the caller:

CREATE TABLE test (col text);
INSERT INTO test VALUES ('123');

CREATE FUNCTION reffunc(refcursor) RETURNS refcursor AS 'BEGIN
OPEN $1 FOR SELECT col FROM test;
RETURN $1;
END;' LANGUAGE plpgsql;

BEGIN;
SELECT reffunc('funccursor');
FETCH ALL IN funccursor;
COMMIT;

The following example uses automatic cursor name generation:

CREATE FUNCTION reffunc2() RETURNS refcursor AS 'DECLARE
ref refcursor;
BEGIN
OPEN ref FOR SELECT col FROM test;
RETURN ref;
END;' LANGUAGE plpgsql;

-- need to be in a transaction to use cursors.
BEGIN;
SELECT reffunc2();

reffunc2
--------------------
<unnamed cursor 1>
(1 row)
FETCH ALL IN "<unnamed cursor 1>";
COMMIT;

The following example shows one way to return multiple cursors from a single function:

CREATE FUNCTION myfunc(refcursor, refcursor) RETURNS SETOF refcursor AS $$
BEGIN
OPEN $1 FOR SELECT * FROM table_1;
RETURN NEXT $1;
OPEN $2 FOR SELECT * FROM table_2;
RETURN NEXT $2;
END;
$$ LANGUAGE plpgsql;
-- need to be in a transaction to use cursors.
BEGIN;

SELECT * FROM myfunc('a', 'b');

FETCH ALL FROM a;
FETCH ALL FROM b;
COMMIT;

41.7.4. Looping Through a Cursor’s Result

There is a variant of the FOR statement that allows iterating through the rows returned by a cursor. The syntax is:

[ <<label>> ]
FOR recordvar IN bound_cursorvar [ ( [ argument_name := ] argument_value [, ...] ) ] LOOP
  statements
END LOOP [ label ];

The cursor variable must have been bound to some query when it was declared, and it cannot be open already. The FOR statement automatically opens the cursor, and it closes the cursor again when the loop exits. A list of actual argument value expressions must appear if and only if the cursor was declared to take arguments. These values will be substituted in the query, in just the same way as during an OPEN (see Section 41.7.2.3).

The variable recordvar is automatically defined as type record and exists only inside the loop (any existing definition of the variable name is ignored within the loop). Each row returned by the cursor is successively assigned to this record variable and the loop body is executed.

41.8. Errors and Messages

41.8.1. Reporting Errors and Messages

Use the RAISE statement to report messages and raise errors.

RAISE [ level ] 'format' [, expression [, ... ] ] [ USING option = expression [, ... ] ];
RAISE [ level ] condition_name [ USING option = expression [, ... ] ];
RAISE [ level ] SQLSTATE 'sqlstate' [ USING option = expression [, ... ] ];
RAISE [ level ] USING option = expression [, ... ];
RAISE ;

The level option specifies the error severity. Allowed levels are DEBUG, LOG, INFO, NOTICE, WARNING, and EXCEPTION, with EXCEPTION being the default. EXCEPTION raises an error (which normally aborts the current transaction); the other levels only generate messages of different priority levels. Whether messages of a particular priority are reported to the client, written to the server log, or both is controlled by the log_min_messages and client_min_messages configuration variables. See Chapter 19 for more information.
After `level` if any, you can write a `format` (which must be a simple string literal, not an expression). The format string specifies the error message text to be reported. The format string can be followed by optional argument expressions to be inserted into the message. Inside the format string, `%` is replaced by the string representation of the next optional argument’s value. Write `%%` to emit a literal `%`. The number of arguments must match the number of `%` placeholders in the format string, or an error is raised during the compilation of the function.

In this example, the value of `v_job_id` will replace the `%` in the string:

```
RAISE NOTICE 'Calling cs_create_job(%)', v_job_id;
```

You can attach additional information to the error report by writing `USING` followed by `option = expression` items. Each `expression` can be any string-valued expression. The allowed `option` key words are:

**MESSAGE**

Sets the error message text. This option can’t be used in the form of `RAISE` that includes a format string before `USING`.

**DETAIL**

Supplies an error detail message.

**HINT**

Supplies a hint message.

**ERRCODE**

Specifies the error code (SQLSTATE) to report, either by condition name, as shown in Appendix A, or directly as a five-character SQLSTATE code.

**COLUMN**

**CONSTRAINT**

**DATATYPE**

**TABLE**

**SCHEMA**

Supplies the name of a related object.

This example will abort the transaction with the given error message and hint:

```
RAISE EXCEPTION 'Nonexistent ID --> %', user_id
    USING HINT = 'Please check your user ID';
```

These two examples show equivalent ways of setting the SQLSTATE:

```
RAISE 'Duplicate user ID: %', user_id USING ERRCODE = 'uniqueViolation';
RAISE 'Duplicate user ID: %', user_id USING ERRCODE = '23505';
```

There is a second `RAISE` syntax in which the main argument is the condition name or SQLSTATE to be reported, for example:

```
RAISE division_by_zero;
```
RAISE SQLSTATE '22012';

In this syntax, USING can be used to supply a custom error message, detail, or hint. Another way to do the earlier example is

RAISE unique_violation USING MESSAGE = 'Duplicate user ID: ' || user_id;

Still another variant is to write RAISE USING or RAISE level USING and put everything else into the USING list.

The last variant of RAISE has no parameters at all. This form can only be used inside a BEGIN block’s EXCEPTION clause; it causes the error currently being handled to be re-thrown.

Note: Before PostgreSQL 9.1, RAISE without parameters was interpreted as re-throwing the error from the block containing the active exception handler. Thus an EXCEPTION clause nested within that handler could not catch it, even if the RAISE was within the nested EXCEPTION clause’s block. This was deemed surprising as well as being incompatible with Oracle’s PL/SQL.

If no condition name nor SQLSTATE is specified in a RAISE EXCEPTION command, the default is to use ERRCODE_RAISE_EXCEPTION (P0001). If no message text is specified, the default is to use the condition name or SQLSTATE as message text.

Note: When specifying an error code by SQLSTATE code, you are not limited to the predefined error codes, but can select any error code consisting of five digits and/or upper-case ASCII letters, other than 00000. It is recommended that you avoid throwing error codes that end in three zeroes, because these are category codes and can only be trapped by trapping the whole category.

41.8.2. Checking Assertions

The ASSERT statement is a convenient shorthand for inserting debugging checks into PL/pgSQL functions.

ASSERT condition [, message ];

The condition is a Boolean expression that is expected to always evaluate to true; if it does, the ASSERT statement does nothing further. If the result is false or null, then an ASSERT_FAILURE exception is raised. (If an error occurs while evaluating the condition, it is reported as a normal error.)

If the optional message is provided, it is an expression whose result (if not null) replaces the default error message text “assertion failed”, should the condition fail. The message expression is not evaluated in the normal case where the assertion succeeds.

Testing of assertions can be enabled or disabled via the configuration parameter plpgsql.check_asserts, which takes a Boolean value; the default is on. If this parameter is off then ASSERT statements do nothing.

Note that ASSERT is meant for detecting program bugs, not for reporting ordinary error conditions. Use the RAISE statement, described above, for that.
41.9. Trigger Procedures

PL/pgSQL can be used to define trigger procedures on data changes or database events. A trigger procedure is created with the `CREATE FUNCTION` command, declaring it as a function with no arguments and a return type of `trigger` (for data change triggers) or `event_trigger` (for database event triggers). Special local variables named `TG_something` are automatically defined to describe the condition that triggered the call.

41.9.1. Triggers on Data Changes

A data change trigger is declared as a function with no arguments and a return type of `trigger`. Note that the function must be declared with no arguments even if it expects to receive some arguments specified in `CREATE TRIGGER` — such arguments are passed via `TG_ARGV`, as described below.

When a PL/pgSQL function is called as a trigger, several special variables are created automatically in the top-level block. They are:

- **NEW**
  - Data type `RECORD`; variable holding the new database row for `INSERT/UPDATE` operations in row-level triggers. This variable is unassigned in statement-level triggers and for `DELETE` operations.

- **OLD**
  - Data type `RECORD`; variable holding the old database row for `UPDATE/DELETE` operations in row-level triggers. This variable is unassigned in statement-level triggers and for `INSERT` operations.

- **TG_NAME**
  - Data type `name`; variable that contains the name of the trigger actually fired.

- **TG_WHEN**
  - Data type `text`; a string of `BEFORE`, `AFTER`, or `INSTEAD OF`, depending on the trigger’s definition.

- **TG_LEVEL**
  - Data type `text`; a string of either `ROW` or `STATEMENT` depending on the trigger’s definition.

- **TG_OP**
  - Data type `text`; a string of `INSERT`, `UPDATE`, `DELETE`, or `TRUNCATE` telling for which operation the trigger was fired.

- **TG_RELID**
  - Data type `oid`; the object ID of the table that caused the trigger invocation.

- **TG_RELNAME**
  - Data type `name`; the name of the table that caused the trigger invocation. This is now deprecated, and could disappear in a future release. Use `TG_TABLE_NAME` instead.

- **TG_TABLE_NAME**
  - Data type `name`; the name of the table that caused the trigger invocation.

- **TG_TABLE_SCHEMA**
  - Data type `name`; the name of the schema of the table that caused the trigger invocation.
TG_NARGS

Data type integer; the number of arguments given to the trigger procedure in the CREATE TRIGGER statement.

TG_ARGV[]

Data type array of text; the arguments from the CREATE TRIGGER statement. The index counts from 0. Invalid indexes (less than 0 or greater than or equal to tg_nargs) result in a null value.

A trigger function must return either NULL or a record/row value having exactly the structure of the table the trigger was fired for.

Row-level triggers fired BEFORE can return null to signal the trigger manager to skip the rest of the operation for this row (i.e., subsequent triggers are not fired, and the INSERT/UPDATE/DELETE does not occur for this row). If a nonnull value is returned then the operation proceeds with that row value. Returning a row value different from the original value of NEW alters the row that will be inserted or updated. Thus, if the trigger function wants the triggering action to succeed normally without altering the row value, NEW (or a value equal thereto) has to be returned. To alter the row to be stored, it is possible to replace single values directly in NEW and return the modified NEW, or to build a complete new record/row to return. In the case of a before-trigger on DELETE, the returned value has no direct effect, but it has to be nonnull to allow the trigger action to proceed. Note that NEW is null in DELETE triggers, so returning that is usually not sensible. The usual idiom in DELETE triggers is to return OLD.

INSTEAD OF triggers (which are always row-level triggers, and may only be used on views) can return null to signal that they did not perform any updates, and that the rest of the operation for this row should be skipped (i.e., subsequent triggers are not fired, and the row is not counted in the rows-affected status for the surrounding INSERT/UPDATE/DELETE). Otherwise a nonnull value should be returned, to signal that the trigger performed the requested operation. For INSERT and UPDATE operations, the return value should be NEW, which the trigger function may modify to support INSERT RETURNING and UPDATE RETURNING (this will also affect the row value passed to any subsequent triggers, or passed to a special EXCLUDED alias reference within an INSERT statement with an ON CONFLICT DO UPDATE clause). For DELETE operations, the return value should be OLD.

The return value of a row-level trigger fired AFTER or a statement-level trigger fired BEFORE or AFTER is always ignored; it might as well be null. However, any of these types of triggers might still abort the entire operation by raising an error.

Example 41-3 shows an example of a trigger procedure in PL/pgSQL.

Example 41-3. A PL/pgSQL Trigger Procedure

This example trigger ensures that any time a row is inserted or updated in the table, the current user name and time are stamped into the row. And it checks that an employee’s name is given and that the salary is a positive value.

```
CREATE TABLE emp (  
    empname text,  
    salary integer,  
    last_date timestamp,  
    last_user text  
);  

CREATE FUNCTION emp_stamp() RETURNS trigger AS $emp_stamp$  
BEGIN  
    -- Check that empname and salary are given  
    IF NEW.empname IS NULL THEN  
        RETURN NULL;  
    END IF;  
    -- Check that salary is positive  
    IF NEW.salary <= 0 THEN  
        RETURN NULL;  
    END IF;  
    -- If both checks pass, return the row  
    RETURN NEW;  
END;  
```

```
```

1157
Chapter 41. PL/pgSQL - SQL Procedural Language

RAISE EXCEPTION 'empname cannot be null';
END IF;

IF NEW.salary IS NULL THEN
    RAISE EXCEPTION '% cannot have null salary', NEW.empname;
END IF;

-- Who works for us when they must pay for it?
IF NEW.salary < 0 THEN
    RAISE EXCEPTION '% cannot have a negative salary', NEW.empname;
END IF;

-- Remember who changed the payroll when
NEW.last_date := current_timestamp;
NEW.last_user := current_user;
RETURN NEW;
END;

$emp_stamp$ LANGUAGE plpgsql;

CREATE TRIGGER emp_stamp BEFORE INSERT OR UPDATE ON emp
FOR EACH ROW EXECUTE PROCEDURE emp_stamp();

Another way to log changes to a table involves creating a new table that holds a row for each insert, update, or delete that occurs. This approach can be thought of as auditing changes to a table. Example 41-4 shows an example of an audit trigger procedure in PL/pgSQL.

Example 41-4. A PL/pgSQL Trigger Procedure For Auditing

This example trigger ensures that any insert, update or delete of a row in the emp table is recorded (i.e., audited) in the emp_audit table. The current time and user name are stamped into the row, together with the type of operation performed on it.

CREATE TABLE emp (
    empname text NOT NULL,
    salary integer
);

CREATE TABLE emp_audit(
    operation char(1) NOT NULL,
    stamp timestamp NOT NULL,
    userid text NOT NULL,
    empname text NOT NULL,
    salary integer
);

CREATE OR REPLACE FUNCTION process_emp_audit() RETURNS TRIGGER AS $emp_audit$
BEGIN
    -- Create a row in emp_audit to reflect the operation performed on emp,
    -- make use of the special variable TG_OP to work out the operation.

    IF (TG_OP = 'DELETE') THEN
        INSERT INTO emp_audit SELECT 'D', now(), user, OLD.*;
        RETURN OLD;
    ELSIF (TG_OP = 'UPDATE') THEN
        INSERT INTO emp_audit SELECT 'U', now(), user, NEW.*;
        RETURN NEW;
    ELSIF (TG_OP = 'INSERT') THEN
        INSERT INTO emp_audit SELECT 'I', now(), user, NEW.*;
        RETURN NEW;
    END IF;
END;

$process_emp_audit$ LANGUAGE plpgsql;

CREATE OR REPLACE FUNCTION process_emp_audit() RETURNS TRIGGER AS $process_emp_audit$
BEGIN
    -- Create a row in emp_audit to reflect the operation performed on emp,
    -- make use of the special variable TG_OP to work out the operation.

    IF (TG_OP = 'DELETE') THEN
        INSERT INTO emp_audit SELECT 'D', now(), user, OLD.*;
        RETURN OLD;
    ELSIF (TG_OP = 'UPDATE') THEN
        INSERT INTO emp_audit SELECT 'U', now(), user, NEW.*;
        RETURN NEW;
    ELSIF (TG_OP = 'INSERT') THEN
        INSERT INTO emp_audit SELECT 'I', now(), user, NEW.*;
        RETURN NEW;
    END IF;
END;

$process_emp_audit$ LANGUAGE plpgsql;

Another way to log changes to a table involves creating a new table that holds a row for each insert, update, or delete that occurs. This approach can be thought of as auditing changes to a table. Example 41-4 shows an example of an audit trigger procedure in PL/pgSQL.

Example 41-4. A PL/pgSQL Trigger Procedure For Auditing

This example trigger ensures that any insert, update or delete of a row in the emp table is recorded (i.e., audited) in the emp_audit table. The current time and user name are stamped into the row, together with the type of operation performed on it.

CREATE TABLE emp (
    empname text NOT NULL,
    salary integer
);

CREATE TABLE emp_audit(
    operation char(1) NOT NULL,
    stamp timestamp NOT NULL,
    userid text NOT NULL,
    empname text NOT NULL,
    salary integer
);

CREATE OR REPLACE FUNCTION process_emp_audit() RETURNS TRIGGER AS $process_emp_audit$
BEGIN
    -- Create a row in emp_audit to reflect the operation performed on emp,
    -- make use of the special variable TG_OP to work out the operation.

    IF (TG_OP = 'DELETE') THEN
        INSERT INTO emp_audit SELECT 'D', now(), user, OLD.*;
        RETURN OLD;
    ELSIF (TG_OP = 'UPDATE') THEN
        INSERT INTO emp_audit SELECT 'U', now(), user, NEW.*;
        RETURN NEW;
    ELSIF (TG_OP = 'INSERT') THEN
        INSERT INTO emp_audit SELECT 'I', now(), user, NEW.*;
        RETURN NEW;
    END IF;
END;

$process_emp_audit$ LANGUAGE plpgsql;
Chapter 41. PL/pgSQL - SQL Procedural Language

`INSERT INTO emp_audit SELECT 'I', now(), user, NEW.*;
RETURN NEW;`

`END IF;
RETURN NULL; -- result is ignored since this is an AFTER trigger
END;`

`$emp_audit$ LANGUAGE plpgsql;`

`CREATE TRIGGER emp_audit
AFTER INSERT OR UPDATE OR DELETE ON emp
    FOR EACH ROW EXECUTE PROCEDURE process_emp_audit();`

A variation of the previous example uses a view joining the main table to the audit table, to show when each entry was last modified. This approach still records the full audit trail of changes to the table, but also presents a simplified view of the audit trail, showing just the last modified timestamp derived from the audit trail for each entry. Example 41-5 shows an example of an audit trigger on a view in PL/pgSQL.

**Example 41-5. A PL/pgSQL View Trigger Procedure For Auditing**

This example uses a trigger on the view to make it updatable, and ensure that any insert, update or delete of a row in the view is recorded (i.e., audited) in the `emp_audit` table. The current time and user name are recorded, together with the type of operation performed, and the view displays the last modified time of each row.

`CREATE TABLE emp {
    empname text PRIMARY KEY,
    salary integer
};`

`CREATE TABLE emp_audit{
    operation char(1) NOT NULL,
    userid text NOT NULL,
    empname text NOT NULL,
    salary integer,
    stamp timestamp NOT NULL
};`

`CREATE VIEW emp_view AS
SELECT e.empname,
    e.salary,
    max(ea.stamp) AS last_updated
FROM emp e
LEFT JOIN emp_audit ea ON ea.empname = e.empname
GROUP BY 1, 2;`

`CREATE OR REPLACE FUNCTION update_emp_view() RETURNS TRIGGER AS $$
BEGIN
    -- Perform the required operation on emp, and create a row in emp_audit
    -- to reflect the change made to emp.

    IF (TG_OP = 'DELETE') THEN
        DELETE FROM emp WHERE empname = OLD.empname;
        IF NOT FOUND THEN RETURN NULL; END IF;
        OLD.last_updated = now();
        INSERT INTO emp_audit VALUES('D', user, OLD.*);
    END IF;
END;$$ LANGUAGE plpgsql;`
RETURN OLD;
ELSEIF (TG_OP = 'UPDATE') THEN
    UPDATE emp SET salary = NEW.salary WHERE empname = OLD.empname;
    IF NOT FOUND THEN RETURN NULL; END IF;

    NEW.last_updated = now();
    INSERT INTO emp_audit VALUES('U', user, NEW.*);
    RETURN NEW;
ELSEIF (TG_OP = 'INSERT') THEN
    INSERT INTO emp VALUES(NEW.empname, NEW.salary);
    NEW.last_updated = now();
    INSERT INTO emp_audit VALUES('I', user, NEW.*);
    RETURN NEW;
END IF;
END;
$$ LANGUAGE plpgsql;

CREATE TRIGGER emp_audit
INSTEAD OF INSERT OR UPDATE OR DELETE ON emp_view
  FOR EACH ROW EXECUTE PROCEDURE update_emp_view();

One use of triggers is to maintain a summary table of another table. The resulting summary can be used in place of the original table for certain queries — often with vastly reduced run times. This technique is commonly used in Data Warehousing, where the tables of measured or observed data (called fact tables) might be extremely large. Example 41-6 shows an example of a trigger procedure in PL/pgSQL that maintains a summary table for a fact table in a data warehouse.

**Example 41-6. A PL/pgSQL Trigger Procedure For Maintaining A Summary Table**

The schema detailed here is partly based on the *Grocery Store* example from *The Data Warehouse Toolkit* by Ralph Kimball.

```sql
-- Main tables - time dimension and sales fact.
--
CREATE TABLE time_dimension (
    time_key integer NOT NULL,
    day_of_week integer NOT NULL,
    day_of_month integer NOT NULL,
    month integer NOT NULL,
    quarter integer NOT NULL,
    year integer NOT NULL
);
CREATE UNIQUE INDEX time_dimension_key ON time_dimension(time_key);

CREATE TABLE sales_fact (
    time_key integer NOT NULL,
    product_key integer NOT NULL,
    store_key integer NOT NULL,
    amount_sold numeric(12,2) NOT NULL,
    units_sold integer NOT NULL,
    amount_cost numeric(12,2) NOT NULL
);
CREATE INDEX sales_fact_time ON sales_fact(time_key);
```

1160
-- Summary table - sales by time.

CREATE TABLE sales_summary_bytime {
  time_key integer NOT NULL,
  amount_sold numeric(15,2) NOT NULL,
  units_sold numeric(12) NOT NULL,
  amount_cost numeric(15,2) NOT NULL
};
CREATE UNIQUE INDEX sales_summary_bytime_key ON sales_summary_bytime(time_key);

-- Function and trigger to amend summarized column(s) on UPDATE, INSERT, DELETE.

CREATE OR REPLACE FUNCTION maint_sales_summary_bytime() RETURNS TRIGGER
AS $maint_sales_summary_bytime$
DECLARE
  delta_time_key integer;
  delta_amount_sold numeric(15,2);
  delta_units_sold numeric(12);
  delta_amount_cost numeric(15,2);
BEGIN
  -- Work out the increment/decrement amount(s).
  IF (TG_OP = 'DELETE') THEN
    delta_time_key = OLD.time_key;
    delta_amount_sold = -1 * OLD.amount_sold;
    delta_units_sold = -1 * OLD.units_sold;
    delta_amount_cost = -1 * OLD.amount_cost;
  ELSIF (TG_OP = 'UPDATE') THEN
    -- forbid updates that change the time_key -
    -- (probably not too onerous, as DELETE + INSERT is how most
    -- changes will be made).
    IF ( OLD.time_key != NEW.time_key) THEN
      RAISE EXCEPTION 'Update of time_key : % -> % not allowed',
                               OLD.time_key, NEW.time_key;
    END IF;
    delta_time_key = OLD.time_key;
    delta_amount_sold = NEW.amount_sold - OLD.amount_sold;
    delta_units_sold = NEW.units_sold - OLD.units_sold;
    delta_amount_cost = NEW.amount_cost - OLD.amount_cost;
  ELSIF (TG_OP = 'INSERT') THEN
    delta_time_key = NEW.time_key;
    delta_amount_sold = NEW.amount_sold;
    delta_units_sold = NEW.units_sold;
    delta_amount_cost = NEW.amount_cost;
  END IF;

  -- Insert or update the summary row with the new values.
  <<insert_update>>

END;
$maint_sales_summary_bytime$;

1161
Chapter 41. PL/pgSQL - SQL Procedural Language

LOOP
UPDATE sales_summary_bytime
    SET amount_sold = amount_sold + delta_amount_sold,
        units_sold = units_sold + delta_units_sold,
        amount_cost = amount_cost + delta_amount_cost
    WHERE time_key = delta_time_key;
EXIT insert_update WHEN found;

BEGIN
    INSERT INTO sales_summary_bytime (
        time_key,
        amount_sold,
        units_sold,
        amount_cost
    )
    VALUES (
        delta_time_key,
        delta_amount_sold,
        delta_units_sold,
        delta_amount_cost
    );
EXIT insert_update;

EXCEPTION
    WHEN UNIQUE_VIOLATION THEN
        -- do nothing
END;
END LOOP insert_update;
RETURN NULL;
END;
$maint_sales_summary_bytime$ LANGUAGE plpgsql;

CREATE TRIGGER maint_sales_summary_bytime
AFTER INSERT OR UPDATE OR DELETE ON sales_fact
FOR EACH ROW EXECUTE PROCEDURE maint_sales_summary_bytime();

INSERT INTO sales_fact VALUES(1,1,1,10,3,15);
INSERT INTO sales_fact VALUES(1,2,1,20,5,35);
INSERT INTO sales_fact VALUES(2,2,1,40,15,135);
INSERT INTO sales_fact VALUES(2,3,1,10,1,13);
SELECT * FROM sales_summary_bytime;
DELETE FROM sales_fact WHERE product_key = 1;
SELECT * FROM sales_summary_bytime;
UPDATE sales_fact SET units_sold = units_sold * 2;
SELECT * FROM sales_summary_bytime;

41.9.2. Triggers on Events

PL/pgSQL can be used to define event triggers. PostgreSQL requires that a procedure that is to be called as an event trigger must be declared as a function with no arguments and a return type of event_trigger.
When a PL/pgSQL function is called as an event trigger, several special variables are created automatically in the top-level block. They are:

**TG_EVENT**
- Data type **text**: a string representing the event the trigger is fired for.

**TG_TAG**
- Data type **text**: variable that contains the command tag for which the trigger is fired.

Example 41-7 shows an example of an event trigger procedure in PL/pgSQL.

**Example 41-7. A PL/pgSQL Event Trigger Procedure**

This example trigger simply raises a NOTICE message each time a supported command is executed.

```sql
CREATE OR REPLACE FUNCTION snitch() RETURNS event_trigger AS $$
BEGIN
  RAISE NOTICE 'snitch: % %', tg_event, tg_tag;
END;
$$ LANGUAGE plpgsql;

CREATE EVENT TRIGGER snitch ON ddl_command_start EXECUTE PROCEDURE snitch();
```

### 41.10. PL/pgSQL Under the Hood

This section discusses some implementation details that are frequently important for PL/pgSQL users to know.

#### 41.10.1. Variable Substitution

SQL statements and expressions within a PL/pgSQL function can refer to variables and parameters of the function. Behind the scenes, PL/pgSQL substitutes query parameters for such references. Parameters will only be substituted in places where a parameter or column reference is syntactically allowed.

As an extreme case, consider this example of poor programming style:

```sql
INSERT INTO foo (foo) VALUES (foo);
```

The first occurrence of `foo` must syntactically be a table name, so it will not be substituted, even if the function has a variable named `foo`. The second occurrence must be the name of a column of the table, so it will not be substituted either. Only the third occurrence is a candidate to be a reference to the function’s variable.

**Note:** PostgreSQL versions before 9.0 would try to substitute the variable in all three cases, leading to syntax errors.

Since the names of variables are syntactically no different from the names of table columns, there can be ambiguity in statements that also refer to tables: is a given name meant to refer to a table column, or a variable? Let’s change the previous example to
INSERT INTO dest (col) SELECT foo + bar FROM src;

Here, dest and src must be table names, and col must be a column of dest, but foo and bar might reasonably be either variables of the function or columns of src.

By default, PL/pgSQL will report an error if a name in a SQL statement could refer to either a variable or a table column. You can fix such a problem by renaming the variable or column, or by qualifying the ambiguous reference, or by telling PL/pgSQL which interpretation to prefer.

The simplest solution is to rename the variable or column. A common coding rule is to use a different naming convention for PL/pgSQL variables than you use for column names. For example, if you consistently name function variables \texttt{v\_something} while none of your column names start with \texttt{v\_}, no conflicts will occur.

Alternatively you can qualify ambiguous references to make them clear. In the above example, src.\texttt{foo} would be an unambiguous reference to the table column. To create an unambiguous reference to a variable, declare it in a labeled block and use the block’s label (see Section 41.2). For example,

\begin{verbatim}
<<block>>>
DECLARE
  foo int;
BEGIN
  foo := ...;
  INSERT INTO dest (col) SELECT block.foo + bar FROM src;

Here block.\texttt{foo} means the variable even if there is a column \texttt{foo} in src. Function parameters, as well as special variables such as \texttt{FOUND}, can be qualified by the function’s name, because they are implicitly declared in an outer block labeled with the function’s name.

Sometimes it is impractical to fix all the ambiguous references in a large body of PL/pgSQL code. In such cases you can specify that PL/pgSQL should resolve ambiguous references as the variable (which is compatible with PL/pgSQL’s behavior before PostgreSQL 9.0), or as the table column (which is compatible with some other systems such as Oracle).

To change this behavior on a system-wide basis, set the configuration parameter \texttt{plpgsql.variable_conflict} to one of \texttt{error}, \texttt{use_variable}, or \texttt{use_column} (where \texttt{error} is the factory default). This parameter affects subsequent compilations of statements in PL/pgSQL functions, but not statements already compiled in the current session. Because changing this setting can cause unexpected changes in the behavior of PL/pgSQL functions, it can only be changed by a superuser.

You can also set the behavior on a function-by-function basis, by inserting one of these special commands at the start of the function text:

\texttt{#variable_conflict error}
\texttt{#variable_conflict use_variable}
\texttt{#variable_conflict use_column}

These commands affect only the function they are written in, and override the setting of \texttt{plpgsql.variable_conflict}. An example is

\begin{verbatim}
CREATE FUNCTION stamp_user(id int, comment text) RETURNS void AS $$
  #variable_conflict use_variable
  DECLARE
    curtime timestamp := now();
  BEGIN
    UPDATE users SET last_modified = curtime, comment = comment
\end{verbatim}
Chapter 41. PL/pgSQL - SQL Procedural Language

WHERE users.id = id;
END;
$$ LANGUAGE plpgsql;

In the UPDATE command, curtime, comment, and id will refer to the function’s variable and parameters whether or not users has columns of those names. Notice that we had to qualify the reference to users.id in the WHERE clause to make it refer to the table column. But we did not have to qualify the reference to comment as a target in the UPDATE list, because syntactically that must be a column of users. We could write the same function without depending on the variable_conflict setting in this way:

CREATE FUNCTION stamp_user(id int, comment text) RETURNS void AS $$
<<fn>>
DECLARE
  curtime timestamp := now();
BEGIN
  UPDATE users SET last_modified = fn.curtime, comment = stamp_user.comment
  WHERE users.id = stamp_user.id;
END;
$$ LANGUAGE plpgsql;

Variable substitution does not happen in the command string given to EXECUTE or one of its variants. If you need to insert a varying value into such a command, do so as part of constructing the string value, or use USING, as illustrated in Section 41.5.4.

Variable substitution currently works only in SELECT, INSERT, UPDATE, and DELETE commands, because the main SQL engine allows query parameters only in these commands. To use a non-constant name or value in other statement types (generically called utility statements), you must construct the utility statement as a string and EXECUTE it.

41.10.2. Plan Caching

The PL/pgSQL interpreter parses the function’s source text and produces an internal binary instruction tree the first time the function is called (within each session). The instruction tree fully translates the PL/pgSQL statement structure, but individual SQL expressions and SQL commands used in the function are not translated immediately.

As each expression and SQL command is first executed in the function, the PL/pgSQL interpreter parses and analyzes the command to create a prepared statement, using the SPI manager’s SPI_prepare function. Subsequent visits to that expression or command reuse the prepared statement. Thus, a function with conditional code paths that are seldom visited will never incur the overhead of analyzing those commands that are never executed within the current session. A disadvantage is that errors in a specific expression or command cannot be detected until that part of the function is reached in execution. (Trivial syntax errors will be detected during the initial parsing pass, but anything deeper will not be detected until execution.)

PL/pgSQL (or more precisely, the SPI manager) can furthermore attempt to cache the execution plan associated with any particular prepared statement. If a cached plan is not used, then a fresh execution plan is generated on each visit to the statement, and the current parameter values (that is, PL/pgSQL variable values) can be used to optimize the selected plan. If the statement has no parameters, or is executed many times, the SPI manager will consider creating a generic plan that is not dependent on specific parameter values, and caching that for re-use. Typically this will happen only if the execution plan is not very sensitive to the values of the PL/pgSQL variables referenced in it. If it is, generating
a plan each time is a net win. See PREPARE for more information about the behavior of prepared statements.

Because PL/pgSQL saves prepared statements and sometimes execution plans in this way, SQL commands that appear directly in a PL/pgSQL function must refer to the same tables and columns on every execution; that is, you cannot use a parameter as the name of a table or column in an SQL command. To get around this restriction, you can construct dynamic commands using the PL/pgSQL EXECUTE statement — at the price of performing new parse analysis and constructing a new execution plan on every execution.

The mutable nature of record variables presents another problem in this connection. When fields of a record variable are used in expressions or statements, the data types of the fields must not change from one call of the function to the next, since each expression will be analyzed using the data type that is present when the expression is first reached. EXECUTE can be used to get around this problem when necessary.

If the same function is used as a trigger for more than one table, PL/pgSQL prepares and caches statements independently for each such table — that is, there is a cache for each trigger function and table combination, not just for each function. This alleviates some of the problems with varying data types; for instance, a trigger function will be able to work successfully with a column named key even if it happens to have different types in different tables.

Likewise, functions having polymorphic argument types have a separate statement cache for each combination of actual argument types they have been invoked for, so that data type differences do not cause unexpected failures.

Statement caching can sometimes have surprising effects on the interpretation of time-sensitive values. For example there is a difference between what these two functions do:

```sql
CREATE FUNCTION logfunc1(logtxt text) RETURNS void AS $$
BEGIN
  INSERT INTO logtable VALUES (logtxt, 'now');
END;
$$ LANGUAGE plpgsql;
```

and:

```sql
CREATE FUNCTION logfunc2(logtxt text) RETURNS void AS $$
DECLARE
  curtime timestamp;
BEGIN
  curtime := 'now';
  INSERT INTO logtable VALUES (logtxt, curtime);
END;
$$ LANGUAGE plpgsql;
```

In the case of logfunc1, the PostgreSQL main parser knows when analyzing the INSERT that the string 'now' should be interpreted as timestamp, because the target column of logtable is of that type. Thus, 'now' will be converted to a timestamp constant when the INSERT is analyzed, and then used in all invocations of logfunc1 during the lifetime of the session. Needless to say, this isn’t what the programmer wanted. A better idea is to use the now() or current_timestamp function.

In the case of logfunc2, the PostgreSQL main parser does not know what type ‘now’ should become and therefore it returns a data value of type text containing the string now. During the ensuing assignment to the local variable curtime, the PL/pgSQL interpreter casts this string to the timestamp type by calling the text_out and timestamp_in functions for the conversion. So, the computed
Chapter 41. PL/pgSQL - SQL Procedural Language

time stamp is updated on each execution as the programmer expects. Even though this happens to work as expected, it’s not terribly efficient, so use of the `now()` function would still be a better idea.

### 41.11. Tips for Developing in PL/pgSQL

One good way to develop in PL/pgSQL is to use the text editor of your choice to create your functions, and in another window, use `psql` to load and test those functions. If you are doing it this way, it is a good idea to write the function using `CREATE OR REPLACE FUNCTION`. That way you can just reload the file to update the function definition. For example:

```sql
CREATE OR REPLACE FUNCTION testfunc(integer) RETURNS integer AS $$
....
$$ LANGUAGE plpgsql;
```

While running `psql`, you can load or reload such a function definition file with:

```
i filename.sql
```

and then immediately issue SQL commands to test the function.

Another good way to develop in PL/pgSQL is with a GUI database access tool that facilitates development in a procedural language. One example of such a tool is pgAdmin, although others exist. These tools often provide convenient features such as escaping single quotes and making it easier to recreate and debug functions.

#### 41.11.1. Handling of Quotation Marks

The code of a PL/pgSQL function is specified in `CREATE FUNCTION` as a string literal. If you write the string literal in the ordinary way with surrounding single quotes, then any single quotes inside the function body must be doubled; likewise any backslashes must be doubled (assuming escape string syntax is used). Doubling quotes is at best tedious, and in more complicated cases the code can become downright incomprehensible, because you can easily find yourself needing half a dozen or more adjacent quote marks. It’s recommended that you instead write the function body as a “dollar-quoted” string literal (see Section 4.1.2.4). In the dollar-quoting approach, you never double any quote marks, but instead take care to choose a different dollar-quoting delimiter for each level of nesting you need. For example, you might write the `CREATE FUNCTION` command as:

```sql
CREATE OR REPLACE FUNCTION testfunc(integer) RETURNS integer AS $PROC$
....
$PROC$ LANGUAGE plpgsql;
```

Within this, you might use quote marks for simple literal strings in SQL commands and `$$` to delimit fragments of SQL commands that you are assembling as strings. If you need to quote text that includes `$`, you could use `$Q`, and so on.

The following chart shows what you have to do when writing quote marks without dollar quoting. It might be useful when translating pre-dollar quoting code into something more comprehensible.
To begin and end the function body, for example:

```
CREATE FUNCTION foo() RETURNS integer AS '.....'
LANGUAGE plpgsql;
```

Anywhere within a single-quoted function body, quote marks must appear in pairs.

For string literals inside the function body, for example:

```
a_output := "Blah";
SELECT * FROM users WHERE f_name="foobar";
```

In the dollar-quoting approach, you’d just write:

```
a_output := 'Blah';
SELECT * FROM users WHERE f_name='foobar';
```

which is exactly what the PL/pgSQL parser would see in either case.

When you need a single quotation mark in a string constant inside the function body, for example:

```
a_output := a_output || " AND name LIKE ""foobar" AND xyz"
```

The value actually appended to `a_output` would be:

```
AND name LIKE 'foobar' AND xyz.
```

In the dollar-quoting approach, you’d write:

```
a_output := a_output || $$ AND name LIKE 'foobar' AND xyz$$
```

being careful that any dollar-quote delimiters around this are not just `$$`.

When a single quotation mark in a string inside the function body is adjacent to the end of that string constant, for example:

```
a_output := a_output || " AND name LIKE ""foobar"
```

The value appended to `a_output` would then be:

```
AND name LIKE 'foobar'.
```

In the dollar-quoting approach, this becomes:

```
a_output := a_output || $$ AND name LIKE 'foobar'$$
```

When you want two single quotation marks in a string constant (which accounts for 8 quotation marks) and this is adjacent to the end of that string constant (2 more). You will probably only need that if you are writing a function that generates other functions, as in Example 41-9. For example:

```
a_output := a_output || " if v_" ||
    referrer_keys.kind || " like """"
    || referrer_keys.key_string || """
    then return """" || referrer_keys.referrer_type
    || """; end if;"
```

The value of `a_output` would then be:

```
if v_... like "..." then return "..."; end if;
```

In the dollar-quoting approach, this becomes:

```
a_output := a_output || $$ if v_$$ || referrer_keys.kind || $$ like '$$
    || referrer_keys.key_string || $$'
    then return '$$ || referrer_keys.referrer_type
    || $$'; end if;$$
```
where we assume we only need to put single quote marks into \texttt{a\_output}, because it will be re-quoted before use.

### 41.11.2. Additional Compile-time Checks

To aid the user in finding instances of simple but common problems before they cause harm, PL/pgSQL provides additional checks. When enabled, depending on the configuration, they can be used to emit either a \texttt{WARNING} or an \texttt{ERROR} during the compilation of a function. A function which has received a \texttt{WARNING} can be executed without producing further messages, so you are advised to test in a separate development environment.

These additional checks are enabled through the configuration variables \texttt{plpgsql.extra_warnings} for warnings and \texttt{plpgsql.extra_errors} for errors. Both can be set either to a comma-separated list of checks, "none" or "all". The default is "none". Currently the list of available checks includes only one:

\texttt{shadowed\_variables}

Checks if a declaration shadows a previously defined variable.

The following example shows the effect of \texttt{plpgsql.extra_warnings} set to \texttt{shadowed\_variables}:

\begin{verbatim}
SET plpgsql.extra_warnings TO 'shadowed\_variables';

CREATE FUNCTION foo(f1 int) RETURNS int AS $$
DECLARE
f1 int;
BEGIN
RETURN f1;
END
$$ LANGUAGE plpgsql;

WARNING: variable "f1" shadows a previously defined variable
LINE 3: f1 int;
   ^
CREATE FUNCTION
\end{verbatim}
a large amount of code that depends on this behavior, setting \texttt{variable\_conflict} may be the best solution.

- In PostgreSQL the function body must be written as a string literal. Therefore you need to use dollar quoting or escape single quotes in the function body. (See Section 41.11.1.)

- Data type names often need translation. For example, in Oracle string values are commonly declared as being of type \texttt{varchar2}, which is a non-SQL-standard type. In PostgreSQL, use type \texttt{varchar} or \texttt{text} instead. Similarly, replace type \texttt{number} with \texttt{numeric}, or use some other numeric data type if there's a more appropriate one.

- Instead of packages, use schemas to organize your functions into groups.

- Since there are no packages, there are no package-level variables either. This is somewhat annoying. You can keep per-session state in temporary tables instead.

- Integer \texttt{FOR} loops with \texttt{REVERSE} work differently: PL/SQL counts down from the second number to the first, while PL/pgSQL counts down from the first number to the second, requiring the loop bounds to be swapped when porting. This incompatibility is unfortunate but is unlikely to be changed. (See Section 41.6.3.5.)

- \texttt{FOR} loops over queries (other than cursors) also work differently: the target variable(s) must have been declared, whereas PL/SQL always declares them implicitly. An advantage of this is that the variable values are still accessible after the loop exits.

- There are various notational differences for the use of cursor variables.

### 41.12.1. Porting Examples

Example 41-8 shows how to port a simple function from PL/SQL to PL/pgSQL.

**Example 41-8. Porting a Simple Function from PL/SQL to PL/pgSQL**

Here is an Oracle PL/SQL function:

```sql
CREATE OR REPLACE FUNCTION cs_fmt_browser_version(v_name varchar2,
                                                 v_version varchar2)
RETURN varchar2 IS
BEGIN
  IF v_version IS NULL THEN
    RETURN v_name;
  END IF;
  RETURN v_name || '/' || v_version;
END;
/
```

Let’s go through this function and see the differences compared to PL/pgSQL:

- The type name \texttt{varchar2} has to be changed to \texttt{varchar} or \texttt{text}. In the examples in this section, we’ll use \texttt{varchar}, but \texttt{text} is often a better choice if you do not need specific string length limits.

- The \texttt{RETURN} key word in the function prototype (not the function body) becomes \texttt{RETURNS} in PostgreSQL. Also, \texttt{IS} becomes \texttt{AS}, and you need to add a \texttt{LANGUAGE} clause because PL/pgSQL is not the only possible function language.
In PostgreSQL, the function body is considered to be a string literal, so you need to use quote marks or dollar quotes around it. This substitutes for the terminating / in the Oracle approach.

The show errors command does not exist in PostgreSQL, and is not needed since errors are reported automatically.

This is how this function would look when ported to PostgreSQL:

```sql
CREATE OR REPLACE FUNCTION cs_fmt_browser_version(v_name varchar, v_version varchar)
RETURNS varchar AS $$
BEGIN
  IF v_version IS NULL THEN
    RETURN v_name;
  END IF;
  RETURN v_name || '/' || v_version;
END;
$$ LANGUAGE plpgsql;
```

Example 41-9 shows how to port a function that creates another function and how to handle the ensuing quoting problems.

**Example 41-9. Porting a Function that Creates Another Function from PL/SQL to PL/pgSQL**

The following procedure grabs rows from a SELECT statement and builds a large function with the results in IF statements, for the sake of efficiency.

This is the Oracle version:

```sql
CREATE OR REPLACE PROCEDURE cs_update_referrer_type_proc IS
  CURSOR referrer_keys IS
    SELECT * FROM cs_referrer_keys
    ORDER BY try_order;
  func_cmd VARCHAR(4000);
BEGIN
  func_cmd := 'CREATE OR REPLACE FUNCTION cs_find_referrer_type(v_host IN VARCHAR2, v_domain IN VARCHAR2, v_url IN VARCHAR2) RETURN VARCHAR2 IS BEGIN';
  FOR referrer_key IN referrer_keys LOOP
    func_cmd := func_cmd || ' IF v_' || referrer_key.kind || ' LIKE '''' || referrer_key.key_string || '''' THEN RETURN '''' || referrer_key.referrer_type || ''''; END IF;'';
  END LOOP;
  func_cmd := func_cmd || ' RETURN NULL; END;';
  EXECUTE IMMEDIATE func_cmd;
END;
/
show errors;
```

Here is how this function would end up in PostgreSQL:

```sql
CREATE OR REPLACE FUNCTION cs_update_referrer_type_proc() RETURNS void AS $func$
DECLARE
```

```sql
```
referrer_keys CURSOR IS
SELECT * FROM cs_referrer_keys
ORDER BY try_order;

func_body text;
func_cmd text;

BEGIN
func_body := 'BEGIN';

FOR referrer_key IN referrer_keys LOOP
  func_body := func_body ||
    ' IF v_' || referrer_key.kind
    || ' LIKE ' || quote_literal(referrer_key.key_string)
    || ' THEN RETURN ' || quote_literal(referrer_key.referrer_type)
    || '; END IF;' ;
END LOOP;

func_body := func_body || ' RETURN NULL; END;';

func_cmd :=
  'CREATE OR REPLACE FUNCTION cs_find_referrer_type(v_host varchar,
    v_domain varchar,
    v_url varchar)
  RETURNS varchar AS
    || quote_literal(func_body)
    || ' LANGUAGE plpgsql;' ;

EXECUTE func_cmd;
END;
$func$ LANGUAGE plpgsql;

Notice how the body of the function is built separately and passed through quote_literal to
double any quote marks in it. This technique is needed because we cannot safely use dollar quoting
for defining the new function: we do not know for sure what strings will be interpolated from the
referrer_key.key_string field. (We are assuming here that referrer_key.kind can be
trusted to always be host, domain, or url, but referrer_key.key_string might be anything,
in particular it might contain dollar signs.) This function is actually an improvement on the
Oracle original, because it will not generate broken code when referrer_key.key_string or
referrer_key.referrer_type contain quote marks.

Example 41-10 shows how to port a function with OUT parameters and string manipulation. Post-
greSQL does not have a built-in instr function, but you can create one using a combination of other
functions. In Section 41.12.3 there is a PL/pgSQL implementation of instr that you can use to make
your porting easier.

Example 41-10. Porting a Procedure With String Manipulation and OUT Parameters from
PL/SQL to PL/pgSQL

The following Oracle PL/SQL procedure is used to parse a URL and return several elements (host, path, and query).

This is the Oracle version:
CREATE OR REPLACE PROCEDURE cs_parse_url(
  v_url IN VARCHAR2,
  v_host OUT VARCHAR2, -- This will be passed back
  v_path OUT VARCHAR2, -- This one too

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Example 41-10. Porting a Procedure With String Manipulation and OUT Parameters from
PL/SQL to PL/pgSQL

The following Oracle PL/SQL procedure is used to parse a URL and return several elements (host, path, and query).

This is the Oracle version:
CREATE OR REPLACE PROCEDURE cs_parse_url(
  v_url IN VARCHAR2,
  v_host OUT VARCHAR2, -- This will be passed back
  v_path OUT VARCHAR2, -- This one too

v_query OUT VARCHAR2) -- And this one
IS
  a_pos1 INTEGER;
  a_pos2 INTEGER;
BEGIN
  v_host := NULL;
  v_path := NULL;
  v_query := NULL;
  a_pos1 := instr(v_url, '//');

  IF a_pos1 = 0 THEN
    RETURN;
  END IF;
  a_pos2 := instr(v_url, '/', a_pos1 + 2);
  IF a_pos2 = 0 THEN
    v_host := substr(v_url, a_pos1 + 2);
    v_path := '/';
    RETURN;
  END IF;
  v_host := substr(v_url, a_pos1 + 2, a_pos2 - a_pos1 - 2);
  a_pos1 := instr(v_url, '?', a_pos2 + 1);

  IF a_pos1 = 0 THEN
    v_path := substr(v_url, a_pos2);
    RETURN;
  END IF;
  v_path := substr(v_url, a_pos2, a_pos1 - a_pos2);
  v_query := substr(v_url, a_pos1 + 1);
END;
/
show errors;

Here is a possible translation into PL/pgSQL:

```sql
CREATE OR REPLACE FUNCTION cs_parse_url(
  v_url IN VARCHAR,
  v_host OUT VARCHAR, -- This will be passed back
  v_path OUT VARCHAR, -- This one too
  v_query OUT VARCHAR) -- And this one
AS $$
DECLARE
  a_pos1 INTEGER;
  a_pos2 INTEGER;
BEGIN
  v_host := NULL;
  v_path := NULL;
  v_query := NULL;
  a_pos1 := instr(v_url, '//');

  IF a_pos1 = 0 THEN
    RETURN;
  END IF;
  a_pos2 := instr(v_url, '/', a_pos1 + 2);
  IF a_pos2 = 0 THEN
    v_host := substr(v_url, a_pos1 + 2);
    v_path := '/';
    RETURN;
  END IF;
  v_host := substr(v_url, a_pos1 + 2, a_pos2 - a_pos1 - 2);
  a_pos1 := instr(v_url, '?', a_pos2 + 1);

  IF a_pos1 = 0 THEN
    v_path := substr(v_url, a_pos2);
    RETURN;
  END IF;
  v_path := substr(v_url, a_pos2, a_pos1 - a_pos2);
  v_query := substr(v_url, a_pos1 + 1);
END;
$$
```

1173
Chapter 41. PL/pgSQL - SQL Procedural Language

RETURN;
END IF;

v_host := substr(v_url, a_pos1 + 2, a_pos2 - a_pos1 - 2);
a_pos1 := instr(v_url, '?', a_pos2 + 1);

IF a_pos1 = 0 THEN
  v_path := substr(v_url, a_pos2);
  RETURN;
END IF;

v_path := substr(v_url, a_pos2, a_pos1 - a_pos2);
v_query := substr(v_url, a_pos1 + 1);
END;
$$ LANGUAGE plpgsql;

This function could be used like this:
SELECT * FROM cs_parse_url('http://foobar.com/query.cgi?baz');

Example 41-11 shows how to port a procedure that uses numerous features that are specific to Oracle.

**Example 41-11. Porting a Procedure from PL/SQL to PL/pgSQL**

The Oracle version:

```sql
CREATE OR REPLACE PROCEDURE cs_create_job(v_job_id IN INTEGER) IS
  a_running_job_count INTEGER;
BEGIN
  LOCK TABLE cs_jobs IN EXCLUSIVE MODE;
  SELECT count(*) INTO a_running_job_count FROM cs_jobs WHERE end_stamp IS NULL;
  IF a_running_job_count > 0 THEN
    COMMIT; -- free lock
    raise_application_error(-20000, 'Unable to create a new job: a job is currently running.);
  END IF;
  DELETE FROM cs_active_job;
  INSERT INTO cs_active_job(job_id) VALUES (v_job_id);
  BEGIN
    INSERT INTO cs_jobs (job_id, start_stamp) VALUES (v_job_id, now());
    EXCEPTION
      WHEN dup_val_on_index THEN NULL; -- don't worry if it already exists
  END;
  COMMIT;
END;
/```

Procedures like this can easily be converted into PostgreSQL functions returning void. This procedure in particular is interesting because it can teach us some things:

1. There is no PRAGMA statement in PostgreSQL.
If you do a `LOCK TABLE` in PL/pgSQL, the lock will not be released until the calling transaction is finished.

You cannot issue `COMMIT` in a PL/pgSQL function. The function is running within some outer transaction and so `COMMIT` would imply terminating the function’s execution. However, in this particular case it is not necessary anyway, because the lock obtained by the `LOCK TABLE` will be released when we raise an error.

This is how we could port this procedure to PL/pgSQL:

```
CREATE OR REPLACE FUNCTION cs_create_job(v_job_id integer) RETURNS void AS $$
DECLARE
  a_running_job_count integer;
BEGIN
  LOCK TABLE cs_jobs IN EXCLUSIVE MODE;
  SELECT count(*) INTO a_running_job_count FROM cs_jobs WHERE end_stamp IS NULL;
  IF a_running_job_count > 0 THEN
    RAISE EXCEPTION 'Unable to create a new job: a job is currently running';
  END IF;
  DELETE FROM cs_active_job;
  INSERT INTO cs_active_job(job_id) VALUES (v_job_id);
  BEGIN
    INSERT INTO cs_jobs (job_id, start_stamp) VALUES (v_job_id, now());
  EXCEPTION
    WHEN unique_violation THEN -- don't worry if it already exists
  END;
END;
$$ LANGUAGE plpgsql;
```

The syntax of `RAISE` is considerably different from Oracle’s statement, although the basic case `RAISE exception_name` works similarly.

The exception names supported by PL/pgSQL are different from Oracle’s. The set of built-in exception names is much larger (see Appendix A). There is not currently a way to declare user-defined exception names, although you can throw user-chosen SQLSTATE values instead.

The main functional difference between this procedure and the Oracle equivalent is that the exclusive lock on the `cs_jobs` table will be held until the calling transaction completes. Also, if the caller later aborts (for example due to an error), the effects of this procedure will be rolled back.

### 41.12.2. Other Things to Watch For

This section explains a few other things to watch for when porting Oracle PL/SQL functions to PostgreSQL.
41.12.2.1. Implicit Rollback after Exceptions

In PL/pgSQL, when an exception is caught by an exception clause, all database changes since the block’s `BEGIN` are automatically rolled back. That is, the behavior is equivalent to what you’d get in Oracle with:

```
BEGIN
    SAVEPOINT s1;
    ... code here ...
EXCEPTION
    WHEN ... THEN
        ROLLBACK TO s1;
        ... code here ...
    WHEN ... THEN
        ROLLBACK TO s1;
        ... code here ...
END;
```

If you are translating an Oracle procedure that uses `SAVEPOINT` and `ROLLBACK TO` in this style, your task is easy: just omit the `SAVEPOINT` and `ROLLBACK TO`. If you have a procedure that uses `SAVEPOINT` and `ROLLBACK TO` in a different way then some actual thought will be required.

41.12.2.2. EXECUTE

The PL/pgSQL version of `EXECUTE` works similarly to the PL/SQL version, but you have to remember to use `quote_literal` and `quote_ident` as described in Section 41.5.4. Constructs of the type `EXECUTE 'SELECT * FROM $1';` will not work reliably unless you use these functions.

41.12.2.3. Optimizing PL/pgSQL Functions

PostgreSQL gives you two function creation modifiers to optimize execution: “volatility” (whether the function always returns the same result when given the same arguments) and “strictness” (whether the function returns null if any argument is null). Consult the CREATE FUNCTION reference page for details.

When making use of these optimization attributes, your `CREATE FUNCTION` statement might look something like this:

```
CREATE FUNCTION foo(...) RETURNS integer AS $$ ...
$$ LANGUAGE plpgsql STRICT IMMUTABLE;
```

41.12.3. Appendix

This section contains the code for a set of Oracle-compatible `instr` functions that you can use to simplify your porting efforts.

```sql
-- instr functions that mimic Oracle's counterpart
-- Syntax: instr(string1, string2 [, n [, m]])
-- where [] denotes optional parameters.
```

1176
-- Search string1, beginning at the nth character, for the mth occurrence
-- of string2. If n is negative, search backwards, starting at the abs(n)’th
-- character from the end of string1.
-- If n is not passed, assume 1 (search starts at first character).
-- If m is not passed, assume 1 (find first occurrence).
-- Returns starting index of string2 in string1, or 0 if string2 is not found.
--
CREATE FUNCTION instr(varchar, varchar) RETURNS integer AS $$_
BEGIN
    RETURN instr($1, $2, 1);
END;
$$ LANGUAGE plpgsql STRICT IMMUTABLE;

CREATE FUNCTION instr(string varchar, string_to_search_for varchar, 
                      beg_index integer)
RETURNS integer AS $$
DECLARE
    pos integer NOT NULL DEFAULT 0;
    temp_str varchar;
    beg integer;
    length integer;
    ss_length integer;
BEGIN
    IF beg_index > 0 THEN
        temp_str := substring(string FROM beg_index);
        pos := position(string_to_search_for IN temp_str);
        IF pos = 0 THEN
            RETURN 0;
        ELSE
            RETURN pos + beg_index - 1;
        END IF;
    ELSIF beg_index < 0 THEN
        ss_length := char_length(string_to_search_for);
        length := char_length(string);
        beg := length + 1 + beg_index;
        WHILE beg > 0 LOOP
            temp_str := substring(string FROM beg FOR ss_length);
            IF string_to_search_for = temp_str THEN
                RETURN beg;
            END IF;
            beg := beg - 1;
        END LOOP;
        RETURN 0;
    ELSE
        RETURN 0;
    END IF;
END;
$$ LANGUAGE plpgsql STRICT IMMUTABLE;
CREATE FUNCTION instr(string varchar, string_to_search_for varchar, 
beg_index integer, occur_index integer) 
RETURNS integer AS $$
DECLARE
    pos integer NOT NULL DEFAULT 0;
    occur_number integer NOT NULL DEFAULT 0;
    temp_str varchar;
    beg integer;
    i integer;
    length integer;
    ss_length integer;
BEGIN
    IF occur_index <= 0 THEN
        RAISE 'argument "%" is out of range', occur_index
            USING ERRCODE = '22003';
    END IF;
    IF beg_index > 0 THEN
        beg := beg_index - 1;
        FOR i IN 1..occur_index LOOP
            temp_str := substring(string FROM beg + 1);
            pos := position(string_to_search_for IN temp_str);
            IF pos = 0 THEN
                RETURN 0;
            END IF;
            beg := beg + pos;
        END LOOP;
        RETURN beg;
    ELSIF beg_index < 0 THEN
        ss_length := char_length(string_to_search_for);
        length := char_length(string);
        beg := length + 1 + beg_index;
        WHILE beg > 0 LOOP
            temp_str := substring(string FROM beg FOR ss_length);
            IF string_to_search_for = temp_str THEN
                occur_number := occur_number + 1;
                IF occur_number = occur_index THEN
                    RETURN beg;
                END IF;
            END IF;
            beg := beg - 1;
        END LOOP;
    ELSE
        RETURN 0;
    END IF;
END;
$$ LANGUAGE plpgsql STRICT IMMUTABLE;
Chapter 42. PL/Tcl - Tcl Procedural Language

PL/Tcl is a loadable procedural language for the PostgreSQL database system that enables the Tcl language\(^1\) to be used to write functions and trigger procedures.

42.1. Overview

PL/Tcl offers most of the capabilities a function writer has in the C language, with a few restrictions, and with the addition of the powerful string processing libraries that are available for Tcl.

One compelling good restriction is that everything is executed from within the safety of the context of a Tcl interpreter. In addition to the limited command set of safe Tcl, only a few commands are available to access the database via SPI and to raise messages via \texttt{elog()}. PL/Tcl provides no way to access internals of the database server or to gain OS-level access under the permissions of the PostgreSQL server process, as a C function can do. Thus, unprivileged database users can be trusted to use this language; it does not give them unlimited authority.

The other notable implementation restriction is that Tcl functions cannot be used to create input/output functions for new data types.

Sometimes it is desirable to write Tcl functions that are not restricted to safe Tcl. For example, one might want a Tcl function that sends email. To handle these cases, there is a variant of PL/Tcl called PL/TclU (for untrusted Tcl). This is exactly the same language except that a full Tcl interpreter is used. If PL/TclU is used, \textit{it must be installed as an untrusted procedural language} so that only database superusers can create functions in it. The writer of a PL/TclU function must take care that the function cannot be used to do anything unwanted, since it will be able to do anything that could be done by a user logged in as the database administrator.

The shared object code for the PL/Tcl and PL/TclU call handlers is automatically built and installed in the PostgreSQL library directory if Tcl support is specified in the configuration step of the installation procedure. To install PL/Tcl and/or PL/TclU in a particular database, use the \texttt{CREATE EXTENSION} command or the \texttt{createlang} program, for example \texttt{createlang pltcl dbname} or \texttt{createlang pltclu dbname}.

42.2. PL/Tcl Functions and Arguments

To create a function in the PL/Tcl language, use the standard \texttt{CREATE FUNCTION} syntax:

```
CREATE FUNCTION funcname (argument-types) RETURNS return-type AS $$
    # PL/Tcl function body
    $$ LANGUAGE pltcl;
```

PL/TclU is the same, except that the language has to be specified as \texttt{pltclu}.

The body of the function is simply a piece of Tcl script. When the function is called, the argument values are passed as variables \$1 ... \$n to the Tcl script. The result is returned from the Tcl code in the usual way, with a \texttt{return} statement.

For example, a function returning the greater of two integer values could be defined as:

```
CREATE FUNCTION tcl_max(integer, integer) RETURNS integer AS $$
    if {$1 > $2} {return $1}
$$
```

\(^1\) http://www.tcl.tk/
return $2
$$ LANGUAGE pltcl STRICT;

Note the clause STRICT, which saves us from having to think about null input values: if a null value is passed, the function will not be called at all, but will just return a null result automatically.

In a nonstrict function, if the actual value of an argument is null, the corresponding $n variable will be set to an empty string. To detect whether a particular argument is null, use the function argisnull. For example, suppose that we wanted tcl_max with one null and one nonnull argument to return the nonnull argument, rather than null:

CREATE FUNCTION tcl_max(integer, integer) RETURNS integer AS $$
if {[argisnull 1]} {
    if {[argisnull 2]} { return_null }
    return $2
} if {[argisnull 2]} { return $1 }
if {{$1 > $2} {return $1}
return $2
$$ LANGUAGE pltcl;

As shown above, to return a null value from a PL/Tcl function, execute return_null. This can be done whether the function is strict or not.

Composite-type arguments are passed to the function as Tcl arrays. The element names of the array are the attribute names of the composite type. If an attribute in the passed row has the null value, it will not appear in the array. Here is an example:

CREATE TABLE employee {
    name text,
    salary integer,
    age integer
};

CREATE FUNCTION overpaid(employee) RETURNS boolean AS $$
if {200000.0 < $1(salary)} {
    return "t"
} if {$1(age) < 30 && 100000.0 < $1(salary)} {
    return "t"
} return "f"
$$ LANGUAGE pltcl;

There is currently no support for returning a composite-type result value, nor for returning sets.

PL/Tcl does not currently have full support for domain types: it treats a domain the same as the underlying scalar type. This means that constraints associated with the domain will not be enforced. This is not an issue for function arguments, but it is a hazard if you declare a PL/Tcl function as returning a domain type.
42.3. Data Values in PL/Tcl

The argument values supplied to a PL/Tcl function’s code are simply the input arguments converted to text form (just as if they had been displayed by a SELECT statement). Conversely, the return command will accept any string that is acceptable input format for the function’s declared return type. So, within the PL/Tcl function, all values are just text strings.

42.4. Global Data in PL/Tcl

Sometimes it is useful to have some global data that is held between two calls to a function or is shared between different functions. This is easily done in PL/Tcl, but there are some restrictions that must be understood.

For security reasons, PL/Tcl executes functions called by any one SQL role in a separate Tcl interpreter for that role. This prevents accidental or malicious interference by one user with the behavior of another user’s PL/Tcl functions. Each such interpreter will have its own values for any “global” Tcl variables. Thus, two PL/Tcl functions will share the same global variables if and only if they are executed by the same SQL role. In an application wherein a single session executes code under multiple SQL roles (via SECURITY DEFINER functions, use of SET ROLE, etc) you may need to take explicit steps to ensure that PL/Tcl functions can share data. To do that, make sure that functions that should communicate are owned by the same user, and mark them SECURITY DEFINER. You must of course take care that such functions can’t be used to do anything unintended.

All PL/Tcl functions used in a session execute in the same Tcl interpreter, which of course is distinct from the interpreter(s) used for PL/Tcl functions. So global data is automatically shared between PL/Tcl functions. This is not considered a security risk because all PL/Tcl functions execute at the same trust level, namely that of a database superuser.

To help protect PL/Tcl functions from unintentionally interfering with each other, a global array is made available to each function via the upvar command. The global name of this variable is the function’s internal name, and the local name is GD. It is recommended that GD be used for persistent private data of a function. Use regular Tcl global variables only for values that you specifically intend to be shared among multiple functions. (Note that the GD arrays are only global within a particular interpreter, so they do not bypass the security restrictions mentioned above.)

An example of using GD appears in the spi_execp example below.

42.5. Database Access from PL/Tcl

The following commands are available to access the database from the body of a PL/Tcl function:

```
spi_exec ?-count n? ?-array name? command ?loop-body?
```

Executes an SQL command given as a string. An error in the command causes an error to be raised. Otherwise, the return value of `spi_exec` is the number of rows processed (selected, inserted, updated, or deleted) by the command, or zero if the command is a utility statement. In addition, if the command is a SELECT statement, the values of the selected columns are placed in Tcl variables as described below.

The optional `-count` value tells `spi_exec` the maximum number of rows to process in the command. The effect of this is comparable to setting up a query as a cursor and then saying `FETCH n`. 
If the command is a `SELECT` statement, the values of the result columns are placed into Tcl variables named after the columns. If the `-array` option is given, the column values are instead stored into elements of the named associative array, with the column names used as array indexes. In addition, the current row number within the result (counting from zero) is stored into the array element named `.tupno`, unless that name is in use as a column name in the result.

If the command is a `SELECT` statement and no `loop-body` script is given, then only the first row of results are stored into Tcl variables or array elements; remaining rows, if any, are ignored. No storing occurs if the query returns no rows. (This case can be detected by checking the result of `spi_exec`.) For example:

```tcl
spi_exec "SELECT count(*) AS cnt FROM pg_proc"
```

will set the Tcl variable `$cnt` to the number of rows in the `pg_proc` system catalog.

If the optional `loop-body` argument is given, it is a piece of Tcl script that is executed once for each row in the query result. (`loop-body` is ignored if the given command is not a `SELECT`.) The values of the current row’s columns are stored into Tcl variables or array elements before each iteration. For example:

```tcl
spi_exec -array C "SELECT * FROM pg_class" {
  elog DEBUG "have table $C(relname)"
}
```

will print a log message for every row of `pg_class`. This feature works similarly to other Tcl looping constructs; in particular `continue` and `break` work in the usual way inside the loop body.

If a column of a query result is null, the target variable for it is “unset” rather than being set.

```tcl
spi_prepare query typelist
```

Prepares and saves a query plan for later execution. The saved plan will be retained for the life of the current session.

The query can use parameters, that is, placeholders for values to be supplied whenever the plan is actually executed. In the query string, refer to parameters by the symbols `$1` ... `$n`. If the query uses parameters, the names of the parameter types must be given as a Tcl list. (Write an empty list for `typelist` if no parameters are used.)

The return value from `spi_prepare` is a query ID to be used in subsequent calls to `spi_execp`. See `spi_execp` for an example.

```tcl
spi_execp ?-count n? ?-array name? ?-nulls string? queryid ?value-list?
?loop-body?
```

Executes a query previously prepared with `spi_prepare`. `queryid` is the ID returned by `spi_prepare`. If the query references parameters, a `value-list` must be supplied. This is a Tcl list of actual values for the parameters. The list must be the same length as the parameter type list previously given to `spi_prepare`. Omit `value-list` if the query has no parameters.

The optional value for `-nulls` is a string of spaces and ‘n’ characters telling `spi_execp` which of the parameters are null values. If given, it must have exactly the same length as the `value-list`. If it is not given, all the parameter values are nonnull.

Except for the way in which the query and its parameters are specified, `spi_execp` works just like `spi_exec`. The `-count`, `-array`, and `loop-body` options are the same, and so is the result value.

Here’s an example of a PL/Tcl function using a prepared plan:

```tcl
CREATE FUNCTION t1_count(integer, integer) RETURNS integer AS $$
  if {![ info exists GD(plan) ]} {
```
Chapter 42. PL/Tcl - Tcl Procedural Language

Chapter 42. PL/Tcl - Tcl Procedural Language

# prepare the saved plan on the first call
set GD(plan) [ spi_prepare "SELECT count(*) AS cnt FROM t1 WHERE num >= \$1 AND num <= \$2" [ list int4 int4 ] ]
spi_execp -count 1 $GD(plan) [ list $1 $2 ]
return $cnt

$$ LANGUAGE pltcl;

We need backslashes inside the query string given to spi_prepare to ensure that the $n markers
will be passed through to spi_prepare as-is, and not replaced by Tcl variable substitution.

spi_lastoid

Returns the OID of the row inserted by the last spi_exec or spi_execp, if the command was
a single-row INSERT and the modified table contained OIDs. (If not, you get zero.)

quote string

Doubles all occurrences of single quote and backslash characters in the given string. This can be
used to safely quote strings that are to be inserted into SQL commands given to spi_exec or
spi_prepare. For example, think about an SQL command string like:

"SELECT '$val' AS ret"

where the Tcl variable val actually contains doesn’t. This would result in the final command string:

SELECT 'doesn’t' AS ret

which would cause a parse error during spi_exec or spi_prepare. To work properly, the
submitted command should contain:

SELECT 'doesn’t' AS ret

which can be formed in PL/Tcl using:

"SELECT '[ quote $val ]' AS ret"

One advantage of spi_execp is that you don’t have to quote parameter values like this, since
the parameters are never parsed as part of an SQL command string.

e log level msg

Emits a log or error message. Possible levels are DEBUG, LOG, INFO, NOTICE, WARNING, ERROR,
and FATAL. ERROR raises an error condition; if this is not trapped by the surrounding Tcl code,
the error propagates out to the calling query, causing the current transaction or subtransaction
to be aborted. This is effectively the same as the Tcl error command. FATAL aborts the trans-
action and causes the current session to shut down. (There is probably no good reason to use
this error level in PL/Tcl functions, but it’s provided for completeness.) The other levels only
generate messages of different priority levels. Whether messages of a particular priority are re-
ported to the client, written to the server log, or both is controlled by the log_min_messages
and client_min_messages configuration variables. See Chapter 19 and Section 42.8 for more
information.

42.6. Trigger Procedures in PL/Tcl

Trigger procedures can be written in PL/Tcl. PostgreSQL requires that a procedure that is to be called
as a trigger must be declared as a function with no arguments and a return type of trigger.

The information from the trigger manager is passed to the procedure body in the following variables:
Chapter 42. PL/Tcl - Tcl Procedural Language

$TG_name
The name of the trigger from the CREATE TRIGGER statement.

$TG_relid
The object ID of the table that caused the trigger procedure to be invoked.

$TG_table_name
The name of the table that caused the trigger procedure to be invoked.

$TG_table_schema
The schema of the table that caused the trigger procedure to be invoked.

$TG_relatts
A Tcl list of the table column names, prefixed with an empty list element. So looking up a column name in the list with Tcl’s lsearch command returns the element’s number starting with 1 for the first column, the same way the columns are customarily numbered in PostgreSQL. (Empty list elements also appear in the positions of columns that have been dropped, so that the attribute numbering is correct for columns to their right.)

$TG_when
The string BEFORE, AFTER, or INSTEAD OF, depending on the type of trigger event.

$TG_level
The string ROW or STATEMENT depending on the type of trigger event.

$TG_op
The string INSERT, UPDATE, DELETE, or TRUNCATE depending on the type of trigger event.

$NEW
An associative array containing the values of the new table row for INSERT or UPDATE actions, or empty for DELETE. The array is indexed by column name. Columns that are null will not appear in the array. This is not set for statement-level triggers.

$OLD
An associative array containing the values of the old table row for UPDATE or DELETE actions, or empty for INSERT. The array is indexed by column name. Columns that are null will not appear in the array. This is not set for statement-level triggers.

$args
A Tcl list of the arguments to the procedure as given in the CREATE TRIGGER statement. These arguments are also accessible as $1 ... $n in the procedure body.

The return value from a trigger procedure can be one of the strings OK or SKIP, or a list of column name/value pairs. If the return value is OK, the operation (INSERT/UPDATE/DELETE) that fired the trigger will proceed normally. SKIP tells the trigger manager to silently suppress the operation for this row. If a list is returned, it tells PL/Tcl to return a modified row to the trigger manager; the contents of the modified row are specified by the column names and values in the list. Any columns not mentioned in the list are set to null. Returning a modified row is only meaningful for row-level BEFORE INSERT or UPDATE triggers, for which the modified row will be inserted instead of the one given in $NEW; or for row-level INSTEAD OF INSERT or UPDATE triggers where the returned row is used as the source data for INSERT RETURNING or UPDATE RETURNING clauses. In row-level BEFORE DELETE or INSTEAD OF DELETE triggers, returning a modified row has the same effect as...
returning OK, that is the operation proceeds. The trigger return value is ignored for all other types of triggers.

**Tip:** The result list can be made from an array representation of the modified tuple with the `array get` Tcl command.

Here’s a little example trigger procedure that forces an integer value in a table to keep track of the number of updates that are performed on the row. For new rows inserted, the value is initialized to 0 and then incremented on every update operation.

```pltcl
CREATE FUNCTION trigfunc_modcount() RETURNS trigger AS $$
switch $TG_op {
INSERT {
    set NEW($1) 0
}
UPDATE {
    set NEW($1) $OLD($1)
    incr NEW($1)
}
default {
    return OK
}
}
return [array get NEW]
$$ LANGUAGE pltcl;
```

CREATE TABLE mytab (num integer, description text, modcnt integer);

CREATE TRIGGER trig_mytab_modcount BEFORE INSERT OR UPDATE ON mytab
FOR EACH ROW EXECUTE PROCEDURE trigfunc_modcount('modcnt');

Notice that the trigger procedure itself does not know the column name; that’s supplied from the trigger arguments. This lets the trigger procedure be reused with different tables.

### 42.7. Event Trigger Procedures in PL/Tcl

Event trigger procedures can be written in PL/Tcl. PostgreSQL requires that a procedure that is to be called as an event trigger must be declared as a function with no arguments and a return type of `event_trigger`.

The information from the trigger manager is passed to the procedure body in the following variables:

- `$TG_event`
  - The name of the event the trigger is fired for.
- `$TG_tag`
  - The command tag for which the trigger is fired.

The return value of the trigger procedure is ignored.

Here’s a little example event trigger procedure that simply raises a `NOTICE` message each time a supported command is executed:
CREATE OR REPLACE FUNCTION tclsnitch() RETURNS event_trigger AS $$
elog NOTICE "tclsnitch: $TG_event $TG_tag"
$$ LANGUAGE pltcl;

CREATE EVENT TRIGGER tcl_a_snitch ON ddl_command_start EXECUTE PROCEDURE tclsnitch();

42.8. Error Handling in PL/Tcl

Tcl code within or called from a PL/Tcl function can raise an error, either by executing some invalid operation or by generating an error using the Tcl error command or PL/Tcl’s elog command. Such errors can be caught within Tcl using the Tcl catch command. If they are not caught but are allowed to propagate out to the top level of execution of the PL/Tcl function, they turn into database errors.

Conversely, database errors that occur within PL/Tcl’s spi_exec, spi_prepare, and spi_execp commands are reported as Tcl errors, so they are catchable by Tcl’s catch command. Again, if they propagate out to the top level without being caught, they turn back into database errors.

Tcl provides an errorCode variable that can represent additional information about an error in a form that is easy for Tcl programs to interpret. The contents are in Tcl list format, and the first word identifies the subsystem or library reporting the error; beyond that the contents are left to the individual subsystem or library. For database errors reported by PL/Tcl commands, the first word is POSTGRES, the second word is the Postgres version number, and additional words are field name/value pairs providing detailed information about the error. Fields SQLSTATE, condition, and message are always supplied (the first two represent the error code and condition name as shown in Appendix A). Fields that may be present include detail, hint, context, schema, table, column, datatype, constraint, statement, cursor_position, filename, lineno, and funcname.

A convenient way to work with PL/Tcl’s errorCode information is to load it into an array, so that the field names become array subscripts. Code for doing that might look like

```tcl
if {[catch { spi_exec $sql_command }]} {
    if {[lindex $::errorCode 0] == "POSTGRES"} {
        array set errorArray $::errorCode
        if {$errorArray(condition) == "undefined_table"} {
            # deal with missing table
        } else {
            # deal with some other type of SQL error
        }
    }
}
```

(The double colons explicitly specify that errorCode is a global variable.)

42.9. Modules and the unknown Command

PL/Tcl has support for autoloading Tcl code when used. It recognizes a special table, pltcl_modules, which is presumed to contain modules of Tcl code. If this table exists, the module unknown is fetched from the table and loaded into the Tcl interpreter immediately before the first execution of a PL/Tcl function in a database session. (This happens separately for each Tcl interpreter, if more than one is used in a session; see Section 42.4.)
While the unknown module could actually contain any initialization script you need, it normally defines a Tcl unknown procedure that is invoked whenever Tcl does not recognize an invoked procedure name. PL/Tcl’s standard version of this procedure tries to find a module in pltcl_modules that will define the required procedure. If one is found, it is loaded into the interpreter, and then execution is allowed to proceed with the originally attempted procedure call. A secondary table pltcl_modfuncs provides an index of which functions are defined by which modules, so that the lookup is reasonably quick.

The PostgreSQL distribution includes support scripts to maintain these tables: pltcl_loadmod, pltcl_listmod, pltcl_delmod, as well as source for the standard unknown module in share/unknown.pltcl. This module must be loaded into each database initially to support the autoloading mechanism.

The tables pltcl_modules and pltcl_modfuncs must be readable by all, but it is wise to make them owned and writable only by the database administrator. As a security precaution, PL/Tcl will ignore pltcl_modules (and thus, not attempt to load the unknown module) unless it is owned by a superuser. But update privileges on this table can be granted to other users, if you trust them sufficiently.

### 42.10. Tcl Procedure Names

In PostgreSQL, the same function name can be used for different function definitions as long as the number of arguments or their types differ. Tcl, however, requires all procedure names to be distinct. PL/Tcl deals with this by making the internal Tcl procedure names contain the object ID of the function from the system table pg_proc as part of their name. Thus, PostgreSQL functions with the same name and different argument types will be different Tcl procedures, too. This is not normally a concern for a PL/Tcl programmer, but it might be visible when debugging.
Chapter 43. PL/Perl - Perl Procedural Language

PL/Perl is a loadable procedural language that enables you to write PostgreSQL functions in the Perl programming language. The main advantage to using PL/Perl is that this allows use, within stored functions, of the manyfold “string munging” operators and functions available for Perl. Parsing complex strings might be easier using Perl than it is with the string functions and control structures provided in PL/pgSQL.

To install PL/Perl in a particular database, use CREATE EXTENSION plperl, or from the shell command line use createlang plperl dbname.

Tip: If a language is installed into template1, all subsequently created databases will have the language installed automatically.

Note: Users of source packages must specially enable the build of PL/Perl during the installation process. (Refer to Chapter 16 for more information.) Users of binary packages might find PL/Perl in a separate subpackage.

43.1. PL/Perl Functions and Arguments

To create a function in the PL/Perl language, use the standard CREATE FUNCTION syntax:

```
CREATE FUNCTION funcname (argument-types) RETURNS return-type AS $$
  # PL/Perl function body
$$ LANGUAGE plperl;
```

The body of the function is ordinary Perl code. In fact, the PL/Perl glue code wraps it inside a Perl subroutine. A PL/Perl function is called in a scalar context, so it can’t return a list. You can return non-scalar values (arrays, records, and sets) by returning a reference, as discussed below.

PL/Perl also supports anonymous code blocks called with the DO statement:

```
DO $$
  # PL/Perl code
$$ LANGUAGE plperl;
```

An anonymous code block receives no arguments, and whatever value it might return is discarded. Otherwise it behaves just like a function.

Note: The use of named nested subroutines is dangerous in Perl, especially if they refer to lexical variables in the enclosing scope. Because a PL/Perl function is wrapped in a subroutine, any named subroutine you place inside one will be nested. In general, it is far safer to create anonymous subroutines which you call via a coderef. For more information, see the entries for Variable "%s" will not stay shared and Variable "%s" is not available in the perl-diag man page, or search the Internet for “perl nested named subroutine”.

1. http://www.perl.org
The syntax of the CREATE FUNCTION command requires the function body to be written as a string constant. It is usually most convenient to use dollar quoting (see Section 4.1.2.4) for the string constant. If you choose to use escape string syntax `E"`, you must double any single quote marks (`'`) and backslashes (`\`) used in the body of the function (see Section 4.1.2.1).

Arguments and results are handled as in any other Perl subroutine: arguments are passed in @_, and a result value is returned with return or as the last expression evaluated in the function.

For example, a function returning the greater of two integer values could be defined as:

```perl
CREATE FUNCTION perl_max (integer, integer) RETURNS integer AS $$
  if ($_[0] > $_[1]) { return $_[0]; }
  return $_[1];
$$ LANGUAGE plperl;
```

**Note:** Arguments will be converted from the database’s encoding to UTF-8 for use inside PL/Perl, and then converted from UTF-8 back to the database encoding upon return.

If an SQL null value is passed to a function, the argument value will appear as “undefined” in Perl. The above function definition will not behave very nicely with null inputs (in fact, it will act as though they are zeroes). We could add STRICT to the function definition to make PostgreSQL do something more reasonable: if a null value is passed, the function will not be called at all, but will just return a null result automatically. Alternatively, we could check for undefined inputs in the function body. For example, suppose that we wanted perl_max with one null and one nonnull argument to return the nonnull argument, rather than a null value:

```perl
CREATE FUNCTION perl_max (integer, integer) RETURNS integer AS $$
  my ($x, $y) = @_; 
  if (not defined $x) {
    return undefined if not defined $y;
    return $y;
  }
  return $x if not defined $y;
  return $x if $x > $y;
  return $y;
$$ LANGUAGE plperl;
```

As shown above, to return an SQL null value from a PL/Perl function, return an undefined value. This can be done whether the function is strict or not.

Anything in a function argument that is not a reference is a string, which is in the standard PostgreSQL external text representation for the relevant data type. In the case of ordinary numeric or text types, Perl will just do the right thing and the programmer will normally not have to worry about it. However, in other cases the argument will need to be converted into a form that is more usable in Perl. For example, the decode_bytea function can be used to convert an argument of type bytea into unescaped binary.

Similarly, values passed back to PostgreSQL must be in the external text representation format. For example, the encode_bytea function can be used to escape binary data for a return value of type bytea.

Perl can return PostgreSQL arrays as references to Perl arrays. Here is an example:

```perl
CREATE OR REPLACE function returns_array()
Perl passes PostgreSQL arrays as a blessed PostgreSQL::InServer::ARRAY object. This object may be treated as an array reference or a string, allowing for backward compatibility with Perl code written for PostgreSQL versions below 9.1 to run. For example:

```perl
CREATE OR REPLACE FUNCTION concat_array_elements(text[]) RETURNS TEXT AS $$
    my $arg = shift;
    my $result = "";
    return undef if (!defined $arg);

    # as an array reference
    for (@$arg) {
        $result .= $_;
    }

    # also works as a string
    $result .= $arg;

    $$ LANGUAGE plperl;
SELECT concat_array_elements(ARRAY['PL','/','Perl']);
```

**Note:** Multidimensional arrays are represented as references to lower-dimensional arrays of references in a way common to every Perl programmer.

Composite-type arguments are passed to the function as references to hashes. The keys of the hash are the attribute names of the composite type. Here is an example:

```sql
CREATE TABLE employee {
    name text,
    basesalary integer,
    bonus integer
};

CREATE FUNCTION empcomp(employee) RETURNS integer AS $$
    my ($emp) = @_;  
    return $emp->{basesalary} + $emp->{bonus};
$$ LANGUAGE plperl;
SELECT name, empcomp(employee.*) FROM employee;
```

A PL/Perl function can return a composite-type result using the same approach: return a reference to a hash that has the required attributes. For example:
CREATE TYPE testrowperl AS (f1 integer, f2 text, f3 text);

CREATE OR REPLACE FUNCTION perl_row() RETURNS testrowperl AS $$
return {f2 => 'hello', f1 => 1, f3 => 'world'};
$$ LANGUAGE plperl;

SELECT * FROM perl_row();

Any columns in the declared result data type that are not present in the hash will be returned as null values.

PL/Perl functions can also return sets of either scalar or composite types. Usually you’ll want to return rows one at a time, both to speed up startup time and to keep from queuing up the entire result set in memory. You can do this with return_next as illustrated below. Note that after the last return_next, you must put either return or (better) return undef.

CREATE OR REPLACE FUNCTION perl_set_int(int) RETURNS SETOF INTEGER AS $$
foreach (0..$_[0]) {
    return_next($_);
}
return undef;
$$ LANGUAGE plperl;

SELECT * FROM perl_set_int(5);

CREATE OR REPLACE FUNCTION perl_set() RETURNS SETOF testrowperl AS $$
return_next({ f1 => 1, f2 => 'Hello', f3 => 'World' });
return_next({ f1 => 2, f2 => 'Hello', f3 => 'PostgreSQL' });
return_next({ f1 => 3, f2 => 'Hello', f3 => 'PL/Perl' });
return undef;
$$ LANGUAGE plperl;

SELECT * FROM perl_set();

For small result sets, you can return a reference to an array that contains either scalars, references to arrays, or references to hashes for simple types, array types, and composite types, respectively. Here are some simple examples of returning the entire result set as an array reference:

CREATE OR REPLACE FUNCTION perl_set_int(int) RETURNS SETOF INTEGER AS $$
return [0..$_[0]];
$$ LANGUAGE plperl;

SELECT * FROM perl_set_int(5);

CREATE OR REPLACE FUNCTION perl_set() RETURNS SETOF testrowperl AS $$
return [
    { f1 => 1, f2 => 'Hello', f3 => 'World' },
    { f1 => 2, f2 => 'Hello', f3 => 'PostgreSQL' },
    { f1 => 3, f2 => 'Hello', f3 => 'PL/Perl' }
];
$$ LANGUAGE plperl;

SELECT * FROM perl_set();

If you wish to use the strict pragma with your code you have a few options. For temporary global use you can SET plperl.use_strict to true. This will affect subsequent compilations of PL/Perl.
functions, but not functions already compiled in the current session. For permanent global use you can set `plperl.use_strict` to true in the `postgresql.conf` file.

For permanent use in specific functions you can simply put:

```perl
use strict;
```

at the top of the function body.

The `feature` pragma is also available to use if your Perl is version 5.10.0 or higher.

### 43.2. Data Values in PL/Perl

The argument values supplied to a PL/Perl function’s code are simply the input arguments converted to text form (just as if they had been displayed by a `SELECT` statement). Conversely, the `return` and `return_next` commands will accept any string that is acceptable input format for the function’s declared return type.

### 43.3. Built-in Functions

#### 43.3.1. Database Access from PL/Perl

Access to the database itself from your Perl function can be done via the following functions:

```perl
spi_exec_query(query [, max-rows])
```

`spi_exec_query` executes an SQL command and returns the entire row set as a reference to an array of hash references. *You should only use this command when you know that the result set will be relatively small.* Here is an example of a query (SELECT command) with the optional maximum number of rows:

```perl
$rv = spi_exec_query('SELECT * FROM my_table', 5);
```

This returns up to 5 rows from the table `my_table`. If `my_table` has a column `my_column`, you can get that value from row `$i` of the result like this:

```perl
$foo = $rv->{rows}[$i]->{my_column};
```

The total number of rows returned from a `SELECT` query can be accessed like this:

```perl
$nrows = $rv->{processed};
```

Here is an example using a different command type:

```perl
$query = "INSERT INTO my_table VALUES (1, 'test')";
$rv = spi_exec_query($query);
```

You can then access the command status (e.g., `SPI_OK_INSERT`) like this:

```perl
$res = $rv->{status};
```

To get the number of rows affected, do:

```perl
$nrows = $rv->{processed};
```

Here is a complete example:

```sql
CREATE TABLE test {
  i int,
  v varchar
};
```
INSERT INTO test (i, v) VALUES (1, 'first line');
INSERT INTO test (i, v) VALUES (2, 'second line');
INSERT INTO test (i, v) VALUES (3, 'third line');
INSERT INTO test (i, v) VALUES (4, 'immortal');

CREATE OR REPLACE FUNCTION test_munge() RETURNS SETOF test AS $$
my $rv = spi_exec_query('select i, v from test;');
my $status = $rv->{status};
my $nrows = $rv->{processed};
foreach my $rn (0 .. $nrows - 1) {
    my $row = $rv->{rows}[$rn];
    $row->{i} += 200 if defined($row->{i});
    $row->{v} =~ tr/A-Za-z/a-zA-Z/ if (defined($row->{v}));
    return_next($row);
}
return undef;
$$ LANGUAGE plperl;

SELECT * FROM test_munge();

spi_query(command)
spi_fetchrow(cursor)
spi_cursor_close(cursor)

spi_query and spi_fetchrow work together as a pair for row sets which might be large, or for cases where you wish to return rows as they arrive. spi_fetchrow works only with spi_query. The following example illustrates how you use them together:

CREATE TYPE foo_type AS (the_num INTEGER, the_text TEXT);

CREATE OR REPLACE FUNCTION lotsa_md5 (INTEGER) RETURNS SETOF foo_type AS $$
use Digest::MD5 qw(md5_hex);
my $file = '/usr/share/dict/words';
my $t = localtime;
elog(NOTICE, "opening file $file at $t" );
open my $fh, '<', $file # ooh, it's a file access!
or elog(ERROR, "cannot open $file for reading: $!" );
my @words = <$fh>;
close $fh;
$t = localtime;
elog(NOTICE, "closed file $file at $t");
chomp(@words);
my $row;
my $sth = spi_query("SELECT * FROM generate_series(1,$_[0]) AS b(a)" );
while (defined ($row = spi_fetchrow($sth))) {
    return_next({
        the_num => $row->{a},
        the_text => md5_hex($words[rand @words])
    });
}
return;
$$ LANGUAGE plperl;

SELECT * from lotsa_md5(500);

Normally, spi_fetchrow should be repeated until it returns undef, indicating that there are no more rows to read. The cursor returned by spi_query is automatically freed.
Chapter 43. PL/Perl - Perl Procedural Language

when spi_fetchrow returns undef. If you do not wish to read all the rows, instead call spi_cursor_close to free the cursor. Failure to do so will result in memory leaks.

spi_prepare(command, argument types)
spi_query_prepared(plan, arguments)
spi_exec_prepared(plan [, attributes], arguments)
spi_freeplan(plan)

spi_prepare, spi_query_prepared, spi_exec_prepared, and spi_freeplan implement the same functionality but for prepared queries. spi_prepare accepts a query string with numbered argument placeholders ($1, $2, etc) and a string list of argument types:

$plan = spi_prepare('SELECT * FROM test WHERE id > $1 AND name = $2', 'INTEGER', 'TEXT');

Once a query plan is prepared by a call to spi_prepare, the plan can be used instead of the string query, either in spi_exec_prepared, where the result is the same as returned by spi_exec_query, or in spi_query_prepared which returns a cursor exactly as spi_query does, which can be later passed to spi_fetchrow. The optional second parameter to spi_exec_prepared is a hash reference of attributes; the only attribute currently supported is limit, which sets the maximum number of rows returned by a query.

The advantage of prepared queries is that it is possible to use one prepared plan for more than one query execution. After the plan is not needed anymore, it can be freed with spi_freeplan:

CREATE OR REPLACE FUNCTION init() RETURNS VOID AS $$
$_SHARED{my_plan} = spi_prepare('SELECT (now() + $1)::date AS now', 'INTERVAL');
$$ LANGUAGE plperl;

CREATE OR REPLACE FUNCTION add_time( INTERVAL ) RETURNS TEXT AS $$
return spi_exec_prepared($_SHARED{my_plan}, $_[0])->{rows}->[0]->{now};
$$ LANGUAGE plperl;

CREATE OR REPLACE FUNCTION done() RETURNS VOID AS $$
spi_freeplan($_SHARED{my_plan});
undef $_SHARED{my_plan};
$$ LANGUAGE plperl;

SELECT init();
SELECT add_time('1 day'), add_time('2 days'), add_time('3 days');
SELECT done();

add_time | add_time | add_time
----------|----------|----------
2005-12-10 | 2005-12-11 | 2005-12-12

Note that the parameter subscript in spi_prepare is defined via $1, $2, $3, etc, so avoid declaring query strings in double quotes that might easily lead to hard-to-catch bugs.

Another example illustrates usage of an optional parameter in spi_exec_prepared:

CREATE TABLE hosts AS SELECT id, ('192.168.1.'||id)::inet AS address
FROM generate_series(1,3) AS id;

CREATE OR REPLACE FUNCTION init_hosts_query() RETURNS VOID AS $$
$_SHARED(plan) = spi_prepare('SELECT * FROM hosts WHERE address << $1', 'inet');
$$ LANGUAGE plperl;
CREATE OR REPLACE FUNCTION query_hosts(inet) RETURNS SETOF hosts AS $$
return spi_exec_prepared(
  $_SHARED{plan},
  {limit => 2},
  $_[0]
)->{rows};
$$ LANGUAGE plperl;

CREATE OR REPLACE FUNCTION release_hosts_query() RETURNS VOID AS $$
spi_freeplan($_SHARED{plan});
undef $_SHARED{plan};
$$ LANGUAGE plperl;

SELECT init_hosts_query();
SELECT query_hosts('192.168.1.0/30');
SELECT release_hosts_query();

query_hosts
-----------------
(1,192.168.1.1)
(2,192.168.1.2)
(2 rows)

43.3.2. Utility Functions in PL/Perl

elog(level, msg)

Emit a log or error message. Possible levels are DEBUG, LOG, INFO, NOTICE, WARNING, and ERROR. ERROR raises an error condition; if this is not trapped by the surrounding Perl code, the error propagates out to the calling query, causing the current transaction or subtransaction to be aborted. This is effectively the same as the Perl die command. The other levels only generate messages of different priority levels. Whether messages of a particular priority are reported to the client, written to the server log, or both is controlled by the log_min_messages and client_min_messages configuration variables. See Chapter 19 for more information.

quote_literal(string)

Return the given string suitably quoted to be used as a string literal in an SQL statement string. Embedded single-quotes and backslashes are properly doubled. Note that quote_literal returns undef on undef input; if the argument might be undef, quote_nullable is often more suitable.

quote_nullable(string)

Return the given string suitably quoted to be used as a string literal in an SQL statement string; or, if the argument is undef, return the unquoted string "NULL". Embedded single-quotes and backslashes are properly doubled.

quote_ident(string)

Return the given string suitably quoted to be used as an identifier in an SQL statement string. Quotes are added only if necessary (i.e., if the string contains non-identifier characters or would be case-folded). Embedded quotes are properly doubled.
decode_bytea(string)

Return the unescaped binary data represented by the contents of the given string, which should be bytea encoded.

encode_bytea(string)

Return the bytea encoded form of the binary data contents of the given string.

encode_array_literal(array)

Returns the contents of the referenced array as a string in array literal format (see Section 8.15.2). Returns the argument value unaltered if it’s not a reference to an array. The delimiter used between elements of the array literal defaults to "", " if a delimiter is not specified or is undef.

encode_typed_literal(value, typename)

Converts a Perl variable to the value of the data type passed as a second argument and returns a string representation of this value. Correctly handles nested arrays and values of composite types.

encode_array_constructor(array)

Returns the contents of the referenced array as a string in array constructor format (see Section 4.2.12). Individual values are quoted using quote_nullable. Returns the argument value, quoted using quote_nullable, if it’s not a reference to an array.

looks_like_number(string)

Returns a true value if the content of the given string looks like a number, according to Perl, returns false otherwise. Returns undef if the argument is undef. Leading and trailing space is ignored. Inf and Infinity are regarded as numbers.

is_array_ref(argument)

Returns a true value if the given argument may be treated as an array reference, that is, if ref of the argument is ARRAY or PostgreSQL::InServer::ARRAY. Returns false otherwise.

---

43.4. Global Values in PL/Perl

You can use the global hash %_SHARED to store data, including code references, between function calls for the lifetime of the current session.

Here is a simple example for shared data:

```perl
CREATE OR REPLACE FUNCTION set_var(name text, val text) RETURNS text AS $$
  if ($_SHARED{$_[0]} = $_[1]) {
    return 'ok';
  } else {
    return "cannot set shared variable $_[0] to $_[1]";
  }
$$ LANGUAGE plperl;

CREATE OR REPLACE FUNCTION get_var(name text) RETURNS text AS $$
  return $_SHARED{$_[0]};
$$ LANGUAGE plperl;

SELECT set_var('sample', 'Hello, PL/Perl! How"s tricks?');
SELECT get_var('sample');
```

1196
Here is a slightly more complicated example using a code reference:

```sql
CREATE OR REPLACE FUNCTION myfuncs() RETURNS void AS $$
$_SHARED{myquote} = sub {
    my $arg = shift;
    $arg =~ s/\('|\\\)/\$1/g;
    return "'\$arg'";
};
$$ LANGUAGE plperl;
SELECT myfuncs(); /* initializes the function */
/* Set up a function that uses the quote function */
CREATE OR REPLACE FUNCTION use_quote(TEXT) RETURNS text AS $$
    my $text_to_quote = shift;
    my $qfunc = $_SHARED{myquote};
    return &$qfunc($text_to_quote);
$$ LANGUAGE plperl;
```

(You could have replaced the above with the one-liner `return $_SHARED{myquote}->{$_[0]};` at the expense of readability.)

For security reasons, PL/Perl executes functions called by any one SQL role in a separate Perl interpreter for that role. This prevents accidental or malicious interference by one user with the behavior of another user’s PL/Perl functions. Each such interpreter has its own value of the `%_SHARED` variable and other global state. Thus, two PL/Perl functions will share the same value of `%_SHARED` if and only if they are executed by the same SQL role. In an application wherein a single session executes code under multiple SQL roles (via `SECURITY DEFINER` functions, use of `SET ROLE`, etc) you may need to take explicit steps to ensure that PL/Perl functions can share data via `%_SHARED`. To do that, make sure that functions that should communicate are owned by the same user, and mark them `SECURITY DEFINER`. You must of course take care that such functions can’t be used to do anything unintended.

### 43.5. Trusted and Untrusted PL/Perl

Normally, PL/Perl is installed as a “trusted” programming language named `plperl`. In this setup, certain Perl operations are disabled to preserve security. In general, the operations that are restricted are those that interact with the environment. This includes file handle operations, `require`, and `use` (for external modules). There is no way to access internals of the database server process or to gain OS-level access with the permissions of the server process, as a C function can do. Thus, any unprivileged database user can be permitted to use this language.

Here is an example of a function that will not work because file system operations are not allowed for security reasons:

```sql
CREATE FUNCTION badfunc() RETURNS integer AS $$
    my $tmpfile = "/tmp/badfile";
    open my $fh, "">", $tmpfile
    or elog(ERROR, qq{could not open the file "$tmpfile": $!});
    print $fh "Testing writing to a file\n";
    close $fh or elog(ERROR, qq{could not close the file "$tmpfile": $!});
    return 1;
$$ LANGUAGE plperl;
```
The creation of this function will fail as its use of a forbidden operation will be caught by the validator. Sometimes it is desirable to write Perl functions that are not restricted. For example, one might want a Perl function that sends mail. To handle these cases, PL/Perl can also be installed as an “untrusted” language (usually called PL/PerlU). In this case the full Perl language is available. When installing the language, the language name `plperlu` will select the untrusted PL/Perl variant.

The writer of a PL/PerlU function must take care that the function cannot be used to do anything unwanted, since it will be able to do anything that could be done by a user logged in as the database administrator. Note that the database system allows only database superusers to create functions in untrusted languages.

If the above function was created by a superuser using the language `plperlu`, execution would succeed.

In the same way, anonymous code blocks written in Perl can use restricted operations if the language is specified as `plperlu` rather than `plperl`, but the caller must be a superuser.

**Note:** While PL/Perl functions run in a separate Perl interpreter for each SQL role, all PL/PerlU functions executed in a given session run in a single Perl interpreter (which is not any of the ones used for PL/Perl functions). This allows PL/PerlU functions to share data freely, but no communication can occur between PL/Perl and PL/PerlU functions.

**Note:** Perl cannot support multiple interpreters within one process unless it was built with the appropriate flags, namely either `usemultiplicity` or `useithreads`. (`usemultiplicity` is preferred unless you actually need to use threads. For more details, see the `perlembed` man page.) If PL/Perl is used with a copy of Perl that was not built this way, then it is only possible to have one Perl interpreter per session, and so any one session can only execute either PL/PerlU functions, or PL/Perl functions that are all called by the same SQL role.

### 43.6. PL/Perl Triggers

PL/Perl can be used to write trigger functions. In a trigger function, the hash reference `$TD` contains information about the current trigger event. `$TD` is a global variable, which gets a separate local value for each invocation of the trigger. The fields of the `$TD` hash reference are:

- `$TD->{new}{foo}`
  - NEW value of column foo
- `$TD->{old}{foo}`
  - OLD value of column foo
- `$TD->{name}`
  - Name of the trigger being called
- `$TD->{event}`
  - Trigger event: INSERT, UPDATE, DELETE, TRUNCATE, or UNKNOWN
- `$TD->{when}`
  - When the trigger was called: BEFORE, AFTER, INSTEAD OF, or UNKNOWN
Row-level triggers can return one of the following:

- `return;`  
  Execute the operation

- "SKIP"  
  Don’t execute the operation

- "MODIFY"  
  Indicates that the `NEW` row was modified by the trigger function

Here is an example of a trigger function, illustrating some of the above:

```sql
CREATE TABLE test (
  i int,
  v varchar
);

CREATE OR REPLACE FUNCTION valid_id() RETURNS trigger AS $$
if ($_TD->{new}{i} >= 100) || ($_TD->{new}{i} <= 0)) {
  return "SKIP";  # skip INSERT/UPDATE command
} elsif ($_TD->{new}{v} ne "immortal") {
  $_TD->{new}{v} .= "(modified by trigger)";
  return "MODIFY";  # modify row and execute INSERT/UPDATE command
} else {
  return;  # execute INSERT/UPDATE command
}
$$ LANGUAGE plperl;

CREATE TRIGGER test_valid_id_trig
  BEFORE INSERT OR UPDATE ON test
```
43.7. PL/Perl Event Triggers

PL/Perl can be used to write event trigger functions. In an event trigger function, the hash reference \$_TD contains information about the current trigger event. \$_TD is a global variable, which gets a separate local value for each invocation of the trigger. The fields of the \$_TD hash reference are:

\$_TD-&gt;{event}

The name of the event the trigger is fired for.

\$_TD-&gt;{tag}

The command tag for which the trigger is fired.

The return value of the trigger procedure is ignored.

Here is an example of an event trigger function, illustrating some of the above:

CREATE OR REPLACE FUNCTION perlsnitch() RETURNS event_trigger AS $$
    elog(NOTICE, "perlsnitch: " . \$_TD-&gt;{event} . " " . \$_TD-&gt;{tag} . " ");
$$ LANGUAGE plperl;

CREATE EVENT TRIGGER perl_a_snitch
    ON ddl_command_start
    EXECUTE PROCEDURE perlsnitch();

43.8. PL/Perl Under the Hood

43.8.1. Configuration

This section lists configuration parameters that affect PL/Perl.

\texttt{plperl.on\_init(string)}

Specifies Perl code to be executed when a Perl interpreter is first initialized, before it is specialized for use by \texttt{plperl} or \texttt{plperlu}. The SPI functions are not available when this code is executed. If the code fails with an error it will abort the initialization of the interpreter and propagate out to the calling query, causing the current transaction or subtransaction to be aborted.

The Perl code is limited to a single string. Longer code can be placed into a module and loaded by the \texttt{on\_init} string. Examples:

\begin{verbatim}
plperl.on_init = 'require "plperlinit.pl"';
plperl.on_init = 'use lib "/my/app"; use MyApp::PgInit;'
\end{verbatim}

Any modules loaded by \texttt{plperl.on\_init}, either directly or indirectly, will be available for use by \texttt{plperl}. This may create a security risk. To see what modules have been loaded you can use:
Initialization will happen in the postmaster if the plperl library is included in shared_preload_libraries, in which case extra consideration should be given to the risk of destabilizing the postmaster. The principal reason for making use of this feature is that Perl modules loaded by plperl.on_init need be loaded only at postmaster start, and will be instantly available without loading overhead in individual database sessions. However, keep in mind that the overhead is avoided only for the first Perl interpreter used by a database session — either PL/PerlU, or PL/Perl for the first SQL role that calls a PL/Perl function. Any additional Perl interpreters created in a database session will have to execute plperl.on_init afresh. Also, on Windows there will be no savings whatsoever from preloading, since the Perl interpreter created in the postmaster process does not propagate to child processes.

This parameter can only be set in the postgresql.conf file or on the server command line.

plperl.on_plperl_init(string)
plperl.on_plperl_init(string)

These parameters specify Perl code to be executed when a Perl interpreter is specialized for plperl or plperl_u respectively. This will happen when a PL/Perl or PL/PerlU function is first executed in a database session, or when an additional interpreter has to be created because the other language is called or a PL/Perl function is called by a new SQL role. This follows any initialization done by plperl.on_init. The SPI functions are not available when this code is executed. The Perl code in plperl.on_plperl_init is executed after "locking down" the interpreter, and thus it can only perform trusted operations.

If the code fails with an error it will abort the initialization and propagate out to the calling query, causing the current transaction or subtransaction to be aborted. Any actions already done within Perl won’t be undone; however, that interpreter won’t be used again. If the language is used again the initialization will be attempted again within a fresh Perl interpreter.

Only superusers can change these settings. Although these settings can be changed within a session, such changes will not affect Perl interpreters that have already been used to execute functions.

plperl.use_strict(boolean)

When set true subsequent compilations of PL/Perl functions will have the strict pragma enabled. This parameter does not affect functions already compiled in the current session.

43.8.2. Limitations and Missing Features

The following features are currently missing from PL/Perl, but they would make welcome contributions.

- PL/Perl functions cannot call each other directly.
- SPI is not yet fully implemented.
- If you are fetching very large data sets using spi_exec_query, you should be aware that these will all go into memory. You can avoid this by using spi_query/spi_fetchrow as illustrated earlier.

A similar problem occurs if a set-returning function passes a large set of rows back to PostgreSQL via return. You can avoid this problem too by instead using return_next for each row returned, as shown previously.
• When a session ends normally, not due to a fatal error, any \texttt{END} blocks that have been defined are executed. Currently no other actions are performed. Specifically, file handles are not automatically flushed and objects are not automatically destroyed.
Chapter 44. PL/Python - Python Procedural Language

The PL/Python procedural language allows PostgreSQL functions to be written in the Python language. 1

To install PL/Python in a particular database, use CREATE EXTENSION plpythonu, or from the shell command line use createlang plpythonu dbname (but see also Section 44.1).

Tip: If a language is installed into template1, all subsequently created databases will have the language installed automatically.

PL/Python is only available as an “untrusted” language, meaning it does not offer any way of restricting what users can do in it and is therefore named plpythonu. A trusted variant plpython might become available in the future if a secure execution mechanism is developed in Python. The writer of a function in untrusted PL/Python must take care that the function cannot be used to do anything unwanted, since it will be able to do anything that could be done by a user logged in as the database administrator. Only superusers can create functions in untrusted languages such as plpythonu.

Note: Users of source packages must specially enable the build of PL/Python during the installation process. (Refer to the installation instructions for more information.) Users of binary packages might find PL/Python in a separate subpackage.

44.1. Python 2 vs. Python 3

PL/Python supports both the Python 2 and Python 3 language variants. (The PostgreSQL installation instructions might contain more precise information about the exact supported minor versions of Python.) Because the Python 2 and Python 3 language variants are incompatible in some important aspects, the following naming and transitioning scheme is used by PL/Python to avoid mixing them:

- The PostgreSQL language named plpython2u implements PL/Python based on the Python 2 language variant.
- The PostgreSQL language named plpython3u implements PL/Python based on the Python 3 language variant.
- The language named plpythonu implements PL/Python based on the default Python language variant, which is currently Python 2. (This default is independent of what any local Python installations might consider to be their “default”, for example, what /usr/bin/python might be.) The default will probably be changed to Python 3 in a distant future release of PostgreSQL, depending on the progress of the migration to Python 3 in the Python community.

This scheme is analogous to the recommendations in PEP 394 2 regarding the naming and transitioning of the python command.

It depends on the build configuration or the installed packages whether PL/Python for Python 2 or Python 3 or both are available.

**Tip:** The built variant depends on which Python version was found during the installation or which version was explicitly set using the `PYTHON` environment variable; see Section 16.4. To make both variants of PL/Python available in one installation, the source tree has to be configured and built twice.

This results in the following usage and migration strategy:

- Existing users and users who are currently not interested in Python 3 use the language name `plpythonu` and don’t have to change anything for the foreseeable future. It is recommended to gradually “future-proof” the code via migration to Python 2.6/2.7 to simplify the eventual migration to Python 3.

  In practice, many PL/Python functions will migrate to Python 3 with few or no changes.

- Users who know that they have heavily Python 2 dependent code and don’t plan to ever change it can make use of the `plpython2u` language name. This will continue to work into the very distant future, until Python 2 support might be completely dropped by PostgreSQL.

- Users who want to dive into Python 3 can use the `plpython3u` language name, which will keep working forever by today’s standards. In the distant future, when Python 3 might become the default, they might like to remove the “3” for aesthetic reasons.

- Daredevils, who want to build a Python-3-only operating system environment, can change the contents of `pg_pltemplate` to make `plpythonu` be equivalent to `plpython3u`, keeping in mind that this would make their installation incompatible with most of the rest of the world.

See also the document What’s New In Python 3.0[^3] for more information about porting to Python 3.

It is not allowed to use PL/Python based on Python 2 and PL/Python based on Python 3 in the same session, because the symbols in the dynamic modules would clash, which could result in crashes of the PostgreSQL server process. There is a check that prevents mixing Python major versions in a session, which will abort the session if a mismatch is detected. It is possible, however, to use both PL/Python variants in the same database, from separate sessions.

### 44.2. PL/Python Functions

Functions in PL/Python are declared via the standard CREATE FUNCTION syntax:

```
CREATE FUNCTION funcname (argument-list)
RETURNS return-type
AS $$
# PL/Python function body
$$ LANGUAGE plpythonu;
```

[^3]: http://docs.python.org/py3k/whatsnew/3.0.html
Chapter 44. PL/Python - Python Procedural Language

The body of a function is simply a Python script. When the function is called, its arguments are passed as elements of the list *args*; named arguments are also passed as ordinary variables to the Python script. Use of named arguments is usually more readable. The result is returned from the Python code in the usual way, with *return* or *yield* (in case of a result-set statement). If you do not provide a return value, Python returns the default *None*. PL/Python translates Python’s *None* into the SQL null value.

For example, a function to return the greater of two integers can be defined as:

```sql
CREATE FUNCTION pymax (a integer, b integer)
RETURNS integer
AS $$
if a > b:
    return a
return b
$$ LANGUAGE plpythonu;
```

The Python code that is given as the body of the function definition is transformed into a Python function. For example, the above results in:

```python
def __plpython_procedure_pymax_23456():
    if a > b:
        return a
    return b
```

assuming that 23456 is the OID assigned to the function by PostgreSQL.

The arguments are set as global variables. Because of the scoping rules of Python, this has the subtle consequence that an argument variable cannot be reassigned inside the function to the value of an expression that involves the variable name itself, unless the variable is redeclared as global in the block. For example, the following won’t work:

```sql
CREATE FUNCTION pystrip(x text)
RETURNS text
AS $$
x = x.strip()  # error
return x
$$ LANGUAGE plpythonu;
```

because assigning to *x* makes *x* a local variable for the entire block, and so the *x* on the right-hand side of the assignment refers to a not-yet-assigned local variable *x*, not the PL/Python function parameter. Using the *global* statement, this can be made to work:

```sql
CREATE FUNCTION pystrip(x text)
RETURNS text
AS $$
global x
x = x.strip()  # ok now
return x
$$ LANGUAGE plpythonu;
```

But it is advisable not to rely on this implementation detail of PL/Python. It is better to treat the function parameters as read-only.
Chapter 44. PL/Python - Python Procedural Language

44.3. Data Values

Generally speaking, the aim of PL/Python is to provide a “natural” mapping between the PostgreSQL and the Python worlds. This informs the data mapping rules described below.

44.3.1. Data Type Mapping

When a PL/Python function is called, its arguments are converted from their PostgreSQL data type to a corresponding Python type:

- PostgreSQL `boolean` is converted to Python `bool`.
- PostgreSQL `smallint` and `int` are converted to Python `int`. PostgreSQL `bigint` and `oid` are converted to `long` in Python 2 and to `int` in Python 3.
- PostgreSQL `real` and `double` are converted to Python `float`.
- PostgreSQL `numeric` is converted to Python `Decimal`. This type is imported from the `cdecimal` package if that is available. Otherwise, `decimal.Decimal` from the standard library will be used. `cdecimal` is significantly faster than `decimal`. In Python 3.3 and up, however, `cdecimal` has been integrated into the standard library under the name `decimal`, so there is no longer any difference.
- PostgreSQL `bytea` is converted to Python `str` in Python 2 and to `bytes` in Python 3. In Python 2, the string should be treated as a byte sequence without any character encoding.
- All other data types, including the PostgreSQL character string types, are converted to a Python `str`. In Python 2, this string will be in the PostgreSQL server encoding; in Python 3, it will be a Unicode string like all strings.
- For nonscalar data types, see below.

When a PL/Python function returns, its return value is converted to the function’s declared PostgreSQL return data type as follows:

- When the PostgreSQL return type is `boolean`, the return value will be evaluated for truth according to the Python rules. That is, 0 and empty string are false, but notably ‘é’ is true.
- When the PostgreSQL return type is `bytea`, the return value will be converted to a string (Python 2) or bytes (Python 3) using the respective Python built-ins, with the result being converted to `bytea`.
- For all other PostgreSQL return types, the return value is converted to a string using the Python built-in `str`, and the result is passed to the input function of the PostgreSQL data type. (If the Python value is a `float`, it is converted using the `repr` built-in instead of `str`, to avoid loss of precision.)

Strings in Python 2 are required to be in the PostgreSQL server encoding when they are passed to PostgreSQL. Strings that are not valid in the current server encoding will raise an error, but not all encoding mismatches can be detected, so garbage data can still result when this is not done correctly. Unicode strings are converted to the correct encoding automatically, so it can be safer and more convenient to use those. In Python 3, all strings are Unicode strings.
- For nonscalar data types, see below.

Note that logical mismatches between the declared PostgreSQL return type and the Python data type of the actual return object are not flagged; the value will be converted in any case.
44.3.2. Null, None

If an SQL null value is passed to a function, the argument value will appear as None in Python. For example, the function definition of `pymax` shown in Section 44.2 will return the wrong answer for null inputs. We could add `STRICT` to the function definition to make PostgreSQL do something more reasonable: if a null value is passed, the function will not be called at all, but will just return a null result automatically. Alternatively, we could check for null inputs in the function body:

```plaintext
CREATE FUNCTION pymax (a integer, b integer) 
RETURNS integer 
AS $$ 
if (a is None) or (b is None): 
    return None 
if a > b: 
    return a 
return b 
$$ LANGUAGE plpythonu;
```

As shown above, to return an SQL null value from a PL/Python function, return the value `None`. This can be done whether the function is strict or not.

44.3.3. Arrays, Lists

SQL array values are passed into PL/Python as a Python list. To return an SQL array value out of a PL/Python function, return a Python sequence, for example a list or tuple:

```plaintext
CREATE FUNCTION return_arr() 
RETURNS int[] 
AS $$ 
return (1, 2, 3, 4, 5) 
$$ LANGUAGE plpythonu;
```

```sql
SELECT return_arr();
```

```
return_arr
-------------
{1,2,3,4,5}
(1 row)
```

Note that in Python, strings are sequences, which can have undesirable effects that might be familiar to Python programmers:

```plaintext
CREATE FUNCTION return_str_arr() 
RETURNS varchar[] 
AS $$ 
return "hello" 
$$ LANGUAGE plpythonu;
```

```sql
SELECT return_str_arr();
```

```
return_str_arr
--------------
{h,e,l,l,o}
(1 row)
```
### 44.3.4. Composite Types

Composite-type arguments are passed to the function as Python mappings. The element names of the mapping are the attribute names of the composite type. If an attribute in the passed row has the null value, it has the value `None` in the mapping. Here is an example:

```sql
CREATE TABLE employee (
    name text,
    salary integer,
    age integer
);

CREATE FUNCTION overpaid (e employee) RETURNS boolean AS $$
    if e["salary"] > 200000:
        return True
    if (e["age"] < 30) and (e["salary"] > 100000):
        return True
    return False
$$ LANGUAGE plpythonu;
```

There are multiple ways to return row or composite types from a Python function. The following examples assume we have:

```sql
CREATE TYPE named_value AS (
    name text,
    value integer
);
```

A composite result can be returned as a:

**Sequence type (a tuple or list, but not a set because it is not indexable)**

Returned sequence objects must have the same number of items as the composite result type has fields. The item with index 0 is assigned to the first field of the composite type, 1 to the second and so on. For example:

```sql
CREATE FUNCTION make_pair (name text, value integer) RETURNS named_value AS $$
    return [ name, value ]
    # or alternatively, as tuple: return ( name, value )
$$ LANGUAGE plpythonu;
```

To return a SQL null for any column, insert `None` at the corresponding position.

**Mapping (dictionary)**

The value for each result type column is retrieved from the mapping with the column name as key. Example:

```sql
CREATE FUNCTION make_pair (name text, value integer) RETURNS named_value AS $$
    return { "name": name, "value": value }
$$ LANGUAGE plpythonu;
```
Any extra dictionary key/value pairs are ignored. Missing keys are treated as errors. To return a SQL null value for any column, insert `None` with the corresponding column name as the key.

Object (any object providing method `__getattr__`)

This works the same as a mapping. Example:

```sql
CREATE FUNCTION make_pair (name text, value integer)
RETURNS named_value
AS $$
class named_value:
    def __init__ (self, n, v):
        self.name = n
        self.value = v
    return named_value(name, value)
# or simply
class nv: pass
nv.name = name
nv.value = value
return nv
$$ LANGUAGE plpythonu;
```

Functions with `OUT` parameters are also supported. For example:

```sql
CREATE FUNCTION multiout_simple(OUT i integer, OUT j integer) AS $$
return (1, 2)
$$ LANGUAGE plpythonu;
```

SELECT * FROM multiout_simple();

### 44.3.5. Set-returning Functions

A PL/Python function can also return sets of scalar or composite types. There are several ways to achieve this because the returned object is internally turned into an iterator. The following examples assume we have composite type:

```sql
CREATE TYPE greeting AS (
    how text,
    who text
);
```

A set result can be returned from a:

Sequence type (tuple, list, set)

```sql
CREATE FUNCTION greet (how text)
RETURNS SETOF greeting
AS $$
# return tuple containing lists as composite types
# all other combinations work also
    return { [ how, "World" ], [ how, "PostgreSQL" ], [ how, "PL/Python" ] }
$$ LANGUAGE plpythonu;
```
Iterator (any object providing __iter__ and next methods)

CREATE FUNCTION greet (how text)
RETURNS SETOF greeting
AS $$
class producer:
    def __init__ (self, how, who):
        self.how = how
        self.who = who
        self.ndx = -1

def __iter__ (self):
    return self

def next (self):
    self.ndx += 1
    if self.ndx == len(self.who):
        raise StopIteration
    return ( self.how, self.who[self.ndx] )

return producer(how, [ "World", "PostgreSQL", "PL/Python" ])
$$ LANGUAGE plpythonu;

Generator (yield)

CREATE FUNCTION greet (how text)
RETURNS SETOF greeting
AS $$
for who in [ "World", "PostgreSQL", "PL/Python" ]:
    yield ( how, who )
$$ LANGUAGE plpythonu;

Warning

Due to Python bug #14831334, some debug versions of Python 2.4 (configured and compiled with option --with-pydebug) are known to crash the PostgreSQL server when using an iterator to return a set result. Unpatched versions of Fedora 4 contain this bug. It does not happen in production versions of Python or on patched versions of Fedora 4.

Set-returning functions with OUT parameters (using RETURNS SETOF record) are also supported. For example:

CREATE FUNCTION multiout_simple_setof(n integer, OUT integer, OUT integer) RETURNS SETOF record
AS $$
return [(1, 2)] * n
$$ LANGUAGE plpythonu;

SELECT * FROM multiout_simple_setof(3);
44.4. Sharing Data

The global dictionary \texttt{SD} is available to store private data between repeated calls to the same function. The global dictionary \texttt{GD} is public data, that is available to all Python functions within a session; use with care.

Each function gets its own execution environment in the Python interpreter, so that global data and function arguments from \texttt{myfunc} are not available to \texttt{myfunc2}. The exception is the data in the \texttt{GD} dictionary, as mentioned above.

44.5. Anonymous Code Blocks

PL/Python also supports anonymous code blocks called with the DO statement:

\begin{verbatim}
DO $$
   # PL/Python code
$$ LANGUAGE plpythonu;
\end{verbatim}

An anonymous code block receives no arguments, and whatever value it might return is discarded. Otherwise it behaves just like a function.

44.6. Trigger Functions

When a function is used as a trigger, the dictionary \texttt{TD} contains trigger-related values:

\begin{itemize}
\item \texttt{TD["event"]} contains the event as a string: \texttt{INSERT}, \texttt{UPDATE}, \texttt{DELETE}, or \texttt{TRUNCATE}.
\item \texttt{TD["when"]} contains one of \texttt{BEFORE}, \texttt{AFTER}, or \texttt{INSTEAD OF}.
\item \texttt{TD["level"]} contains \texttt{ROW} or \texttt{STATEMENT}.
\item \texttt{TD["new"]}
\item \texttt{TD["old"]}
\end{itemize}

For a row-level trigger, one or both of these fields contain the respective trigger rows, depending on the trigger event.

\begin{itemize}
\item \texttt{TD["name"]} contains the trigger name.
\item \texttt{TD["table_name"]}
\item \texttt{TD["table_schema"]}
\item \texttt{TD["relid"]}
\end{itemize}

contains the OID of the table on which the trigger occurred.
TD["args"]

If the CREATE TRIGGER command included arguments, they are available in TD["args"][0] to TD["args"][n-1].

If TD["when"] is BEFORE or INSTEAD OF and TD["level"] is ROW, you can return None or "OK" from the Python function to indicate the row is unmodified, "SKIP" to abort the event, or if TD["event"] is INSERT or UPDATE you can return "MODIFY" to indicate you’ve modified the new row. Otherwise the return value is ignored.

44.7. Database Access

The PL/Python language module automatically imports a Python module called plpy. The functions and constants in this module are available to you in the Python code as plpy.foo.

44.7.1. Database Access Functions

The plpy module provides several functions to execute database commands:

plpy.execute(query [, max-rows])

Calling plpy.execute with a query string and an optional row limit argument causes that query to be run and the result to be returned in a result object.

The result object emulates a list or dictionary object. The result object can be accessed by row number and column name. For example:

rv = plpy.execute("SELECT * FROM my_table", 5)
returns up to 5 rows from my_table. If my_table has a column my_column, it would be accessed as:

foo = rv[i]["my_column"]
The number of rows returned can be obtained using the built-in len function.

The result object has these additional methods:

nrows()
Returns the number of rows processed by the command. Note that this is not necessarily the same as the number of rows returned. For example, an UPDATE command will set this value but won’t return any rows (unless RETURNING is used).

status()
The SPI_execute() return value.

colnames()
coltypes()
coltypmods()

Return a list of column names, list of column type OIDs, and list of type-specific type modifiers for the columns, respectively.

These methods raise an exception when called on a result object from a command that did not produce a result set, e.g., UPDATE without RETURNING, or DROP TABLE. But it is OK to use these methods on a result set containing zero rows.
The standard `__str__` method is defined so that it is possible for example to debug query execution results using `plpy.debug(rv)`.

The result object can be modified. Note that calling `plpy.execute` will cause the entire result set to be read into memory. Only use that function when you are sure that the result set will be relatively small. If you don’t want to risk excessive memory usage when fetching large results, use `plpy.cursor` rather than `plpy.execute`.

```python
plpy.prepare(query [, argtypes])
plpy.execute(plan [, arguments [, max-rows]])
```

`plpy.prepare` prepares the execution plan for a query. It is called with a query string and a list of parameter types, if you have parameter references in the query. For example:

```python
plan = plpy.prepare("SELECT last_name FROM my_users WHERE first_name = $1", ["text"])
text is the type of the variable you will be passing for $1. The second argument is optional if you don’t want to pass any parameters to the query.
```

After preparing a statement, you use a variant of the function `plpy.execute` to run it:

```python
rv = plpy.execute(plan, ["name"], 5)
```

Pass the plan as the first argument (instead of the query string), and a list of values to substitute into the query as the second argument. The second argument is optional if the query does not expect any parameters. The third argument is the optional row limit as before.

Query parameters and result row fields are converted between PostgreSQL and Python data types as described in Section 44.3.

When you prepare a plan using the PL/Python module it is automatically saved. Read the SPI documentation (Chapter 45) for a description of what this means. In order to make effective use of this across function calls one needs to use one of the persistent storage dictionaries `SD` or `GD` (see Section 44.4). For example:

```sql
CREATE FUNCTION usesavedplan() RETURNS trigger AS $$
if "plan" in SD:
    plan = SD["plan"]
else:
    plan = plpy.prepare("SELECT 1")
SD["plan"] = plan
# rest of function
$$ LANGUAGE plpythonu;
```

```python
plpy.cursor(query)
plpy.cursor(plan [, arguments])
```

The `plpy.cursor` function accepts the same arguments as `plpy.execute` (except for the row limit) and returns a cursor object, which allows you to process large result sets in smaller chunks. As with `plpy.execute`, either a query string or a plan object along with a list of arguments can be used.

The cursor object provides a `fetch` method that accepts an integer parameter and returns a result object. Each time you call `fetch`, the returned object will contain the next batch of rows, never larger than the parameter value. Once all rows are exhausted, `fetch` starts returning an empty result object. Cursor objects also provide an iterator interface, yielding one row at a time until

5. [http://docs.python.org/library/stdtypes.html#iterator-types](http://docs.python.org/library/stdtypes.html#iterator-types)
all rows are exhausted. Data fetched that way is not returned as result objects, but rather as
dictionaries, each dictionary corresponding to a single result row.

An example of two ways of processing data from a large table is:

```
CREATE FUNCTION count_odd_iterator() RETURNS integer AS $$
odd = 0
for row in plpy.cursor("select num from largetable"):  
    if row['num'] % 2:
        odd += 1
return odd
$$ LANGUAGE plpythonu;

CREATE FUNCTION count_odd_fetch(batch_size integer) RETURNS integer AS $$
odd = 0
cursor = plpy.cursor("select num from largetable")
while True:
    rows = cursor.fetch(batch_size)
    if not rows:
        break
    for row in rows:
        if row['num'] % 2:
            odd += 1
return odd
$$ LANGUAGE plpythonu;

CREATE FUNCTION count_odd_prepared() RETURNS integer AS $$
odd = 0
plan = plpy.prepare("select num from largetable where num % $1 <> 0", ["integer"]) 
rows = list(plpy.cursor(plan, [2]))
return len(rows)
$$ LANGUAGE plpythonu;
```

Cursors are automatically disposed of. But if you want to explicitly release all resources held by
a cursor, use the `close` method. Once closed, a cursor cannot be fetched from anymore.

Tip: Do not confuse objects created by `plpy.cursor` with DB-API cursors as defined by
the Python Database API specification\(^6\). They don't have anything in common except for the
name.

### 44.7.2. Trapping Errors

Functions accessing the database might encounter errors, which will cause them to abort and raise
an exception. Both `plpy.execute` and `plpy.prepare` can raise an instance of a subclass of
`plpy.SPILError`, which by default will terminate the function. This error can be handled just like
any other Python exception, by using the `try/except` construct. For example:

```
CREATE FUNCTION try_adding_joe() RETURNS text AS $$
try:
    plpy.execute("INSERT INTO users(username) VALUES ('joe')")
$$
```

---

\(^6\) [http://www.python.org/dev/peps/pep-0249/](http://www.python.org/dev/peps/pep-0249/)
except plpy.SPIError:
    return "something went wrong"
else:
    return "Joe added"
$$ LANGUAGE plpythonu;

The actual class of the exception being raised corresponds to the specific condition that caused the error. Refer to Table A-1 for a list of possible conditions. The module `plpy.spiexceptions` defines an exception class for each PostgreSQL condition, deriving their names from the condition name. For instance, `division_by_zero` becomes `DivisionByZero`, `unique_violation` becomes `UniqueViolation`, `fdw_error` becomes `FdwError`, and so on. Each of these exception classes inherits from `SPIError`. This separation makes it easier to handle specific errors, for instance:

```sql
CREATE FUNCTION insert_fraction(numerator int, denominator int) RETURNS text AS $$
from plpy import spiexceptions
try:
    plan = plpy.prepare("INSERT INTO fractions (frac) VALUES ($1 / $2)", ["int", "int"])
    plpy.execute(plan, [numerator, denominator])
except spiexceptions.DivisionByZero:
    return "denominator cannot equal zero"
except spiexceptions.UniqueViolation:
    return "already have that fraction"
except plpy.SPIError, e:
    return "other error, SQLSTATE %s" % e.sqlstate
else:
    return "fraction inserted"
$$ LANGUAGE plpythonu;
```

Note that because all exceptions from the `plpy.spiexceptions` module inherit from `SPIError`, an `except` clause handling it will catch any database access error.

As an alternative way of handling different error conditions, you can catch the `SPIError` exception and determine the specific error condition inside the `except` block by looking at the `sqlstate` attribute of the exception object. This attribute is a string value containing the “SQLSTATE” error code. This approach provides approximately the same functionality.

### 44.8. Explicit Subtransactions

Recovering from errors caused by database access as described in Section 44.7.2 can lead to an undesirable situation where some operations succeed before one of them fails, and after recovering from that error the data is left in an inconsistent state. PL/Python offers a solution to this problem in the form of explicit subtransactions.

#### 44.8.1. Subtransaction Context Managers

Consider a function that implements a transfer between two accounts:

```sql
CREATE FUNCTION transfer_funds() RETURNS void AS $$
try:
    plpy.execute("UPDATE accounts SET balance = balance - 100 WHERE account_name = 'joe'
    plpy.execute("UPDATE accounts SET balance = balance + 100 WHERE account_name = 'mary'
```
except plpy_SPILError, e:
    result = "error transferring funds: %s" % e.args
else:
    result = "funds transferred correctly"
plan = plpy.prepare("INSERT INTO operations (result) VALUES ($1)", ["text"])
plpy.execute(plan, [result])
$$ LANGUAGE plpythonu;

If the second UPDATE statement results in an exception being raised, this function will report the error, but the result of the first UPDATE will nevertheless be committed. In other words, the funds will be withdrawn from Joe’s account, but will not be transferred to Mary’s account.

To avoid such issues, you can wrap your plpy.execute calls in an explicit subtransaction. The plpy module provides a helper object to manage explicit subtransactions that gets created with the plpy.subtransaction() function. Objects created by this function implement the context manager interface. Using explicit subtransactions we can rewrite our function as:

CREATE FUNCTION transfer_funds2() RETURNS void AS $$
try:
    with plpy.subtransaction():
        plpy.execute("UPDATE accounts SET balance = balance - 100 WHERE account_name = 'joe'")
        plpy.execute("UPDATE accounts SET balance = balance + 100 WHERE account_name = 'mary'")
except plpy_SPILError, e:
    result = "error transferring funds: %s" % e.args
else:
    result = "funds transferred correctly"
plan = plpy.prepare("INSERT INTO operations (result) VALUES ($1)", ["text"])
plpy.execute(plan, [result])
$$ LANGUAGE plpythonu;

Note that the use of try/catch is still required. Otherwise the exception would propagate to the top of the Python stack and would cause the whole function to abort with a PostgreSQL error, so that the operations table would not have any row inserted into it. The subtransaction context manager does not trap errors, it only assures that all database operations executed inside its scope will be atomically committed or rolled back. A rollback of the subtransaction block occurs on any kind of exception exit, not only ones caused by errors originating from database access. A regular Python exception raised inside an explicit subtransaction block would also cause the subtransaction to be rolled back.

### 44.8.2. Older Python Versions

Context managers syntax using the with keyword is available by default in Python 2.6. If using PL/Python with an older Python version, it is still possible to use explicit subtransactions, although not as transparently. You can call the subtransaction manager’s __enter__ and __exit__ functions using the enter and exit convenience aliases. The example function that transfers funds could be written as:

CREATE FUNCTION transfer_funds_old() RETURNS void AS $$
try:
    subxact = plpy.subtransaction()
    subxact.enter()
    try:
        plpy.execute("UPDATE accounts SET balance = balance - 100 WHERE account_name = 'joe'")
        plpy.execute("UPDATE accounts SET balance = balance + 100 WHERE account_name = 'mary'")
    except plpy_SPILError, e:
        result = "error transferring funds: %s" % e.args
    else:
        result = "funds transferred correctly"
    plan = plpy.prepare("INSERT INTO operations (result) VALUES ($1)", ["text"])
    plpy.execute(plan, [result])
$$ LANGUAGE plpythonu;

7. [http://docs.python.org/library/stdtypes.html#context-manager-types](http://docs.python.org/library/stdtypes.html#context-manager-types)
except:
    import sys
    subxact.exit(*sys.exc_info())
    raise
else:
    subxact.exit(None, None, None)
except plpy.SPIError, e:
    result = "error transferring funds: %s" % e.args
else:
    result = "funds transferred correctly"

plan = plpy.prepare("INSERT INTO operations (result) VALUES ($1)", ["text"])
plpy.execute(plan, [result])
$$ LANGUAGE plpythonu;

Note: Although context managers were implemented in Python 2.5, to use the with syntax in that version you need to use a future statement. Because of implementation details, however, you cannot use future statements in PL/Python functions.

44.9. Utility Functions

The plpy module also provides the functions

plpy.debug(msg, **kwargs)
plpy.log(msg, **kwargs)
plpy.info(msg, **kwargs)
plpy.notice(msg, **kwargs)
plpy.warning(msg, **kwargs)
plpy.error(msg, **kwargs)
plpy.fatal(msg, **kwargs)

plpy.error and plpy.fatal actually raise a Python exception which, if uncaught, propagates out to the calling query, causing the current transaction or subtransaction to be aborted. raise plpy.Error(msg) and raise plpy.Fatal(msg) are equivalent to calling plpy.error(msg) and plpy.fatal(msg), respectively but the raise form does not allow passing keyword arguments. The other functions only generate messages of different priority levels. Whether messages of a particular priority are reported to the client, written to the server log, or both is controlled by the log_min_messages and client_min_messages configuration variables. See Chapter 19 for more information.

The msg argument is given as a positional argument. For backward compatibility, more than one positional argument can be given. In that case, the string representation of the tuple of positional arguments becomes the message reported to the client.

The following keyword-only arguments are accepted:

detail

---

The string representation of the objects passed as keyword-only arguments is used to enrich the messages reported to the client. For example:

```
CREATE FUNCTION raise_custom_exception() RETURNS void AS $$
    plpy.error("custom exception message",
               detail="some info about exception",
               hint="hint for users")
$$ LANGUAGE plpythonu;
```

```
=# SELECT raise_custom_exception();
ERROR: plpy.Error: custom exception message
DETAIL: some info about exception
HINT: hint for users
CONTEXT: Traceback (most recent call last):
    PL/Python function "raise_custom_exception", line 4, in <module>
        hint="hint for users")
PL/Python function "raise_custom_exception"
```

Another set of utility functions are `plpy.quote_literal(string)`, `plpy.quote_nullable(string)`, and `plpy.quote_ident(string)`. They are equivalent to the built-in quoting functions described in Section 9.4. They are useful when constructing ad-hoc queries. A PL/Python equivalent of dynamic SQL from Example 41-1 would be:

```
plpy.execute("UPDATE tbl SET %s = %s WHERE key = %s" %
              (plpy.quote_ident(colname),
               plpy.quote_nullable(newvalue),
               plpy.quote_literal(keyvalue)))
```

### 44.10. Environment Variables

Some of the environment variables that are accepted by the Python interpreter can also be used to affect PL/Python behavior. They would need to be set in the environment of the main PostgreSQL server process, for example in a start script. The available environment variables depend on the version of Python; see the Python documentation for details. At the time of this writing, the following environment variables have an affect on PL/Python, assuming an adequate Python version:

- PYTHONHOME
- PYTHONPATH
- PYTHON2K
- PYTHONOPTIMIZE
Chapter 44. PL/Python - Python Procedural Language

- PYTHONDEBUG
- PYTHONVERBOSE
- PYTHONCASEOK
- PYTHONDONTWRITEBYTECODE
- PYTHONIOENCODING
- PYTHONUSERBASE
- PYTHONHASHSEED

(It appears to be a Python implementation detail beyond the control of PL/Python that some of the environment variables listed on the python man page are only effective in a command-line interpreter and not an embedded Python interpreter.)
Chapter 45. Server Programming Interface

The Server Programming Interface (SPI) gives writers of user-defined C functions the ability to run SQL commands inside their functions. SPI is a set of interface functions to simplify access to the parser, planner, and executor. SPI also does some memory management.

Note: The available procedural languages provide various means to execute SQL commands from procedures. Most of these facilities are based on SPI, so this documentation might be of use for users of those languages as well.

To avoid misunderstanding we’ll use the term “function” when we speak of SPI interface functions and “procedure” for a user-defined C-function that is using SPI.

Note that if a command invoked via SPI fails, then control will not be returned to your procedure. Rather, the transaction or subtransaction in which your procedure executes will be rolled back. (This might seem surprising given that the SPI functions mostly have documented error-return conventions. Those conventions only apply for errors detected within the SPI functions themselves, however.) It is possible to recover control after an error by establishing your own subtransaction surrounding SPI calls that might fail.

SPI functions return a nonnegative result on success (either via a returned integer value or in the global variable SPI_result, as described below). On error, a negative result or NULL will be returned.

Source code files that use SPI must include the header file executor/spi.h.

45.1. Interface Functions

SPI_connect

Name

SPI_connect — connect a procedure to the SPI manager

Synopsis

int SPI_connect(void)

Description

SPI_connect opens a connection from a procedure invocation to the SPI manager. You must call this function if you want to execute commands through SPI. Some utility SPI functions can be called from unconnected procedures.

If your procedure is already connected, SPI_connect will return the error code SPI_ERROR_CONNECT. This could happen if a procedure that has called SPI_connect directly calls another procedure that calls SPI_connect. While recursive calls to the SPI manager are
permitted when an SQL command called through SPI invokes another function that uses SPI, directly
nested calls to SPI_connect and SPI_finish are forbidden. (But see SPI_push and SPI_pop.)

Return Value

SPI_OK_CONNECT
  on success
SPI_ERROR_CONNECT
  on error
SPI_finish

Name

SPI_finish — disconnect a procedure from the SPI manager

Synopsis

int SPI_finish(void)

Description

SPI_finish closes an existing connection to the SPI manager. You must call this function after completing the SPI operations needed during your procedure’s current invocation. You do not need to worry about making this happen, however, if you abort the transaction via \texttt{elog(ERROR)}. In that case SPI will clean itself up automatically.

If \texttt{SPI_finish} is called without having a valid connection, it will return \texttt{SPI_ERROR_UNCONNECTED}. There is no fundamental problem with this; it means that the SPI manager has nothing to do.

Return Value

\texttt{SPI_OK_FINISH}

if properly disconnected

\texttt{SPI_ERROR_UNCONNECTED}

if called from an unconnected procedure
SPI_push

Name

SPI_push — push SPI stack to allow recursive SPI usage

Synopsis

void SPI_push(void)

Description

SPI_push should be called before executing another procedure that might itself wish to use SPI. After SPI_push, SPI is no longer in a “connected” state, and SPI function calls will be rejected unless a fresh SPI_connect is done. This ensures a clean separation between your procedure’s SPI state and that of another procedure you call. After the other procedure returns, call SPI_pop to restore access to your own SPI state.

Note that SPI_execute and related functions automatically do the equivalent of SPI_push before passing control back to the SQL execution engine, so it is not necessary for you to worry about this when using those functions. Only when you are directly calling arbitrary code that might contain SPI_connect calls do you need to issue SPI_push and SPI_pop.
SPI_pop

Name

SPI_pop — pop SPI stack to return from recursive SPI usage

Synopsis

void SPI_pop(void)

Description

SPI_pop pops the previous environment from the SPI call stack. See SPI_push.
**Name**

SPI_execute — execute a command

**Synopsis**

```c
int SPI_execute(const char * command, bool read_only, long count)
```

**Description**

SPI_execute executes the specified SQL command for `count` rows. If `read_only` is true, the command must be read-only, and execution overhead is somewhat reduced.

This function can only be called from a connected procedure.

If `count` is zero then the command is executed for all rows that it applies to. If `count` is greater than zero, then no more than `count` rows will be retrieved; execution stops when the count is reached, much like adding a LIMIT clause to the query. For example,

```
SPI_execute("SELECT * FROM foo", true, 5);
```

will retrieve at most 5 rows from the table. Note that such a limit is only effective when the command actually returns rows. For example,

```
SPI_execute("INSERT INTO foo SELECT * FROM bar", false, 5);
```

inserts all rows from `bar`, ignoring the `count` parameter. However, with

```
SPI_execute("INSERT INTO foo SELECT * FROM bar RETURNING *", false, 5);
```

at most 5 rows would be inserted, since execution would stop after the fifth RETURNING result row is retrieved.

You can pass multiple commands in one string; SPI_execute returns the result for the command executed last. The `count` limit applies to each command separately (even though only the last result will actually be returned). The limit is not applied to any hidden commands generated by rules.

When `read_only` is false, SPI_execute increments the command counter and computes a new snapshot before executing each command in the string. The snapshot does not actually change if the current transaction isolation level is SERIALIZABLE or REPEATABLE READ, but in READ COMMITTED mode the snapshot update allows each command to see the results of newly committed transactions from other sessions. This is essential for consistent behavior when the commands are modifying the database.

When `read_only` is true, SPI_execute does not update either the snapshot or the command counter, and it allows only plain SELECT commands to appear in the command string. The commands are executed using the snapshot previously established for the surrounding query. This execution mode is somewhat faster than the read/write mode due to eliminating per-command overhead. It also allows genuinely stable functions to be built: since successive executions will all use the same snapshot, there will be no change in the results.
SPI_execute

It is generally unwise to mix read-only and read-write commands within a single function using SPI; that could result in very confusing behavior, since the read-only queries would not see the results of any database updates done by the read-write queries.

The actual number of rows for which the (last) command was executed is returned in the global variable SPI_processed. If the return value of the function is SPI_OK_SELECT, SPI_OK_INSERT_RETURNING, SPI_OK_DELETE_RETURNING, or SPI_OK_UPDATE_RETURNING, then you can use the global pointer SPITupleTable *SPI_tuptable to access the result rows. Some utility commands (such as EXPLAIN) also return row sets, and SPI_tuptable will contain the result in these cases too. Some utility commands (COPY, CREATE TABLE AS) don’t return a row set, so SPI_tuptable is NULL, but they still return the number of rows processed in SPI_processed.

The structure SPITupleTable is defined thus:

typdef struct
{
    MemoryContext tuptabcxt;    /* memory context of result table */
    uint64 allocated;           /* number of allocated vals */
    uint64 free;                /* number of free vals */
    TupleDesc tupdesc;          /* row descriptor */
    HeapTuple *vals;            /* rows */
} SPITupleTable;

vals is an array of pointers to rows. (The number of valid entries is given by SPI_processed.) tupdesc is a row descriptor which you can pass to SPI functions dealing with rows. tuptabcxt, allocated, and free are internal fields not intended for use by SPI callers.

SPI_finish frees all SPITupleTable*s allocated during the current procedure. You can free a particular result table earlier, if you are done with it, by calling SPI_freetuptable.

Arguments

const char * command
    string containing command to execute
bool read_only
    true for read-only execution
long count
    maximum number of rows to return, or 0 for no limit

Return Value

If the execution of the command was successful then one of the following (nonnegative) values will be returned:

SPI_OK_SELECT
    if a SELECT (but not SELECT INTO) was executed
SPI_OK_SELINTO
    if a SELECT INTO was executed
SPI_OK_INSERT
    if an INSERT was executed
SPI_OK_DELETE
    if a DELETE was executed
SPI_OK_UPDATE
    if an UPDATE was executed
SPI_OK_INSERT_RETURNING
    if an INSERT RETURNING was executed
SPI_OK_DELETE_RETURNING
    if a DELETE RETURNING was executed
SPI_OK_UPDATE_RETURNING
    if an UPDATE RETURNING was executed
SPI_OK_UTILITY
    if a utility command (e.g., CREATE TABLE) was executed
SPI_OK_REWRITTEN
    if the command was rewritten into another kind of command (e.g., UPDATE became an INSERT) by a rule.

On error, one of the following negative values is returned:

SPI_ERROR_ARGUMENT
    if command is NULL or count is less than 0
SPI_ERROR_COPY
    if COPY TO stdout or COPY FROM stdin was attempted
SPI_ERROR_TRANSACTION
    if a transaction manipulation command was attempted (BEGIN, COMMIT, ROLLBACK, SAVEPOINT, PREPARE TRANSACTION, COMMIT PREPARED, ROLLBACK PREPARED, or any variant thereof)
SPI_ERROR_OPUNKNOWN
    if the command type is unknown (shouldn’t happen)
SPI_ERROR_UNCONNECTED
    if called from an unconnected procedure

Notes

All SPI query-execution functions set both SPI_processed and SPI_tuptable (just the pointer, not the contents of the structure). Save these two global variables into local procedure variables if you need to access the result table of SPI_execute or another query-execution function across later calls.
SPI_exec

Name
SPI_exec — execute a read/write command

Synopsis
int SPI_exec(const char * command, long count)

Description
SPI_exec is the same as SPI_execute, with the latter’s read_only parameter always taken as false.

Arguments
const char * command
    string containing command to execute
long count
    maximum number of rows to return, or 0 for no limit

Return Value
See SPI_execute.
SPI_execute_with_args

Name

SPI_execute_with_args — execute a command with out-of-line parameters

Synopsis

int SPI_execute_with_args(const char *command,
    int nargs, Oid *argtypes,
    Datum *values, const char *nulls,
    bool read_only, long count)

Description

SPI_execute_with_args executes a command that might include references to externally supplied parameters. The command text refers to a parameter as \$n, and the call specifies data types and values for each such symbol. read_only and count have the same interpretation as in SPI_execute.

The main advantage of this routine compared to SPI_execute is that data values can be inserted into the command without tedious quoting/escaping, and thus with much less risk of SQL-injection attacks.

Similar results can be achieved with SPI_prepare followed by SPI_execute_plan; however, when using this function the query plan is always customized to the specific parameter values provided. For one-time query execution, this function should be preferred. If the same command is to be executed with many different parameters, either method might be faster, depending on the cost of re-planning versus the benefit of custom plans.

Arguments

const char * command
    command string

int nargs
    number of input parameters (\$1, \$2, etc.)

Oid * argtypes
    an array of length nargs, containing the OIDs of the data types of the parameters

Datum * values
    an array of length nargs, containing the actual parameter values

const char * nulls
    an array of length nargs, describing which parameters are null

If nulls is NULL then SPI_execute_with_args assumes that no parameters are null. Otherwise, each entry of the nulls array should be ‘ ’ if the corresponding parameter value is non-null, or ‘n’ if the corresponding parameter value is null. (In the latter case, the actual value
in the corresponding values entry doesn’t matter.) Note that nulls is not a text string, just an 
array: it does not need a ‘\0’ terminator.

bool read_only
    true for read-only execution
long count
    maximum number of rows to return, or 0 for no limit

Return Value

The return value is the same as for SPI_execute.

SPI_processed and SPI_tuptable are set as in SPI_execute if successful.
**SPI_prepare**

**Name**

SPI_prepare — prepare a statement, without executing it yet

**Synopsis**

SPIPlanPtr SPI_prepare(const char * command, int nargs, Oid * argtypes)

**Description**

SPI_prepare creates and returns a prepared statement for the specified command, but doesn’t execute the command. The prepared statement can later be executed repeatedly using SPI_execute_plan.

When the same or a similar command is to be executed repeatedly, it is generally advantageous to perform parse analysis only once, and might furthermore be advantageous to re-use an execution plan for the command. SPI_prepare converts a command string into a prepared statement that encapsulates the results of parse analysis. The prepared statement also provides a place for caching an execution plan if it is found that generating a custom plan for each execution is not helpful.

A prepared command can be generalized by writing parameters ($1, $2, etc.) in place of what would be constants in a normal command. The actual values of the parameters are then specified when SPI_execute_plan is called. This allows the prepared command to be used over a wider range of situations than would be possible without parameters.

The statement returned by SPI_prepare can be used only in the current invocation of the procedure, since SPI_finish frees memory allocated for such a statement. But the statement can be saved for longer using the functions SPI_keepplan or SPI_saveplan.

**Arguments**

const char * command
  command string

int nargs
  number of input parameters ($1, $2, etc.)

Oid * argtypes
  pointer to an array containing the OIDs of the data types of the parameters

**Return Value**

SPI_prepare returns a non-null pointer to an SPIPlan, which is an opaque struct representing a prepared statement. On error, NULL will be returned, and SPI_result will be set to one of the same error codes used by SPI_execute, except that it is set to SPI_ERROR_ARGUMENT if command is NULL, or if nargs is less than 0, or if nargs is greater than 0 and argtypes is NULL.
Notes

If no parameters are defined, a generic plan will be created at the first use of \texttt{SPI_execute_plan}, and used for all subsequent executions as well. If there are parameters, the first few uses of \texttt{SPI_execute_plan} will generate custom plans that are specific to the supplied parameter values. After enough uses of the same prepared statement, \texttt{SPI_execute_plan} will build a generic plan, and if that is not too much more expensive than the custom plans, it will start using the generic plan instead of re-planning each time. If this default behavior is unsuitable, you can alter it by passing the \texttt{CURSOR\_OPT\_GENERIC\_PLAN} or \texttt{CURSOR\_OPT\_CUSTOM\_PLAN} flag to \texttt{SPI_prepare_cursor}, to force use of generic or custom plans respectively.

Although the main point of a prepared statement is to avoid repeated parse analysis and planning of the statement, PostgreSQL will force re-analysis and re-planning of the statement before using it whenever database objects used in the statement have undergone definitional (DDL) changes since the previous use of the prepared statement. Also, if the value of search_path changes from one use to the next, the statement will be re-parsed using the new search_path. (This latter behavior is new as of PostgreSQL 9.3.) See PREPARE for more information about the behavior of prepared statements.

This function should only be called from a connected procedure.

\texttt{SPIPlanPtr} is declared as a pointer to an opaque struct type in \texttt{spi.h}. It is unwise to try to access its contents directly, as that makes your code much more likely to break in future revisions of PostgreSQL.

The name \texttt{SPIPlanPtr} is somewhat historical, since the data structure no longer necessarily contains an execution plan.
SPI_prepare_cursor

Name
SPI_prepare_cursor — prepare a statement, without executing it yet

Synopsis
SPIPlanPtr SPI_prepare_cursor(const char * command, int nargs,
    Oid * argtypes, int cursorOptions)

Description
SPI_prepare_cursor is identical to SPI_prepare, except that it also allows specification
of the planner’s “cursor options” parameter. This is a bit mask having the values shown in
nodes/parsenodes.h for the options field of DeclareCursorStmt. SPI_prepare always
takes the cursor options as zero.

Arguments
const char * command
    command string
int nargs
    number of input parameters ($1, $2, etc.)
Oid * argtypes
    pointer to an array containing the OIDs of the data types of the parameters
int cursorOptions
    integer bit mask of cursor options; zero produces default behavior

Return Value
SPI_prepare_cursor has the same return conventions as SPI_prepare.

Notes
Useful bits to set in cursorOptions include CURSOR_OPT_SCROLL, CURSOR_OPT_NO_SCROLL,
CURSOR_OPT_FAST_PLAN, CURSOR_OPT_GENERIC_PLAN, and CURSOR_OPT_CUSTOM_PLAN. Note
in particular that CURSOR_OPT_HOLD is ignored.
**SPI_prepare_params**

**Name**

SPI_prepare_params — prepare a statement, without executing it yet

**Synopsis**

SPIPlanPtr SPI_prepare_params(const char *command,
                               ParserSetupHook parserSetup,
                               void *parserSetupArg,
                               int cursorOptions)

**Description**

SPI_prepare_params creates and returns a prepared statement for the specified command, but doesn’t execute the command. This function is equivalent to SPI_prepare_cursor, with the addition that the caller can specify parser hook functions to control the parsing of external parameter references.

**Arguments**

const char * command  
command string

ParserSetupHook parserSetup
Parser hook setup function

void * parserSetupArg
pass-through argument for parserSetup

int cursorOptions
integer bit mask of cursor options; zero produces default behavior

**Return Value**

SPI_prepare_params has the same return conventions as SPI_prepare.
SPI_getargcount

**Name**

SPI_getargcount — return the number of arguments needed by a statement prepared by SPI_prepare

**Synopsis**

```c
int SPI_getargcount(SPIPlanPtr plan)
```

**Description**

SPI_getargcount returns the number of arguments needed to execute a statement prepared by SPI_prepare.

**Arguments**

SPIPlanPtr plan

prepared statement (returned by SPI_prepare)

**Return Value**

The count of expected arguments for the plan. If the plan is NULL or invalid, SPI_result is set to SPI_ERROR_ARGUMENT and -1 is returned.
SPI_getargtypeid

**Name**

SPI_getargtypeid — return the data type OID for an argument of a statement prepared by SPI_prepare

**Synopsis**

Oid SPI_getargtypeid(SPIPlanPtr plan, int argIndex)

**Description**

SPI_getargtypeid returns the OID representing the type for the argIndex’th argument of a statement prepared by SPI_prepare. First argument is at index zero.

**Arguments**

SPIPlanPtr plan

prevented statement (returned by SPI_prepare)

int argIndex

zero based index of the argument

**Return Value**

The type OID of the argument at the given index. If the plan is NULL or invalid, or argIndex is less than 0 or not less than the number of arguments declared for the plan, SPI_result is set to SPI_ERROR_ARGUMENT and InvalidOid is returned.
SPI_is_cursor_plan

Name

SPI_is_cursor_plan — return true if a statement prepared by SPIPrepare can be used with SPI_cursor_open

Synopsis

bool SPI_is_cursor_plan(SPIPlanPtr plan)

Description

SPI_is_cursor_plan returns true if a statement prepared by SPI_prepare can be passed as an argument to SPI_cursor_open, or false if that is not the case. The criteria are that the plan represents one single command and that this command returns tuples to the caller; for example, SELECT is allowed unless it contains an INTO clause, and UPDATE is allowed only if it contains a RETURNING clause.

Arguments

SPIPlanPtr plan

prepared statement (returned by SPI_prepare)

Return Value

true or false to indicate if the plan can produce a cursor or not, with SPI_result set to zero. If it is not possible to determine the answer (for example, if the plan is NULL or invalid, or if called when not connected to SPI), then SPI_result is set to a suitable error code and false is returned.
**SPI_execute_plan**

**Name**
SPI_execute_plan — execute a statement prepared by SPI_prepare

**Synopsis**

```c
int SPI_execute_plan(SPIPlanPtr plan, Datum * values, const char * nulls,
                     bool read_only, long count)
```

**Description**

SPI_execute_plan executes a statement prepared by SPI_prepare or one of its siblings. `read_only` and `count` have the same interpretation as in SPI_execute.

**Arguments**

SPIPlanPtr plan
  prepared statement (returned by SPI_prepare)

Datum * values
  An array of actual parameter values. Must have same length as the statement’s number of arguments.

const char * nulls
  An array describing which parameters are null. Must have same length as the statement’s number of arguments.
  If `nulls` is NULL then SPI_execute_plan assumes that no parameters are null. Otherwise, each entry of the `nulls` array should be ‘ ’ if the corresponding parameter value is non-null, or ‘n’ if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding `values` entry doesn’t matter.) Note that `nulls` is not a text string, just an array: it does not need a ‘\0’ terminator.

bool read_only
  true for read-only execution

long count
  maximum number of rows to return, or 0 for no limit
Return Value

The return value is the same as for `SPI_execute`, with the following additional possible error (negative) results:

- **SPI_ERROR_ARGUMENT**
  - if `plan` is NULL or invalid, or `count` is less than 0
- **SPI_ERROR_PARAM**
  - if `values` is NULL and `plan` was prepared with some parameters

`SPI_processed` and `SPI_tuptable` are set as in `SPI_execute` if successful.
SPI_execute_plan_with_paramlist

Name

SPI_execute_plan_with_paramlist — execute a statement prepared by SPI_prepare

Synopsis

int SPI_execute_plan_with_paramlist(SPIPlanPtr plan,
                                  ParamListInfo params,
                                  bool read_only,
                                  long count)

Description

SPI_execute_plan_with_paramlist executes a statement prepared by SPI_prepare. This function is equivalent to SPI_execute_plan except that information about the parameter values to be passed to the query is presented differently. The ParamListInfo representation can be convenient for passing down values that are already available in that format. It also supports use of dynamic parameter sets via hook functions specified in ParamListInfo.

Arguments

SPIPlanPtr plan
    prepared statement (returned by SPI_prepare)
ParamListInfo params
    data structure containing parameter types and values; NULL if none
bool read_only
    true for read-only execution
long count
    maximum number of rows to return, or 0 for no limit

Return Value

The return value is the same as for SPI_execute_plan. SPI_processed and SPI_tuptable are set as in SPI_execute_plan if successful.
SPI_execep

Name
SPI_execep — execute a statement in read/write mode

Synopsis
int SPI_execep(SPIPlanPtr plan, Datum * values, const char * nulls, long count)

Description
SPI_execep is the same as SPI_execute_plan, with the latter’s read_only parameter always taken as false.

Arguments
SPIPlanPtr plan
    prepared statement (returned by SPI_prepare)
Datum * values
    An array of actual parameter values. Must have same length as the statement’s number of arguments.
const char * nulls
    An array describing which parameters are null. Must have same length as the statement’s number of arguments.
    If nulls is NULL then SPI_execep assumes that no parameters are null. Otherwise, each entry of the nulls array should be ‘ ’ if the corresponding parameter value is non-null, or ‘n’ if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding values entry doesn’t matter.) Note that nulls is not a text string, just an array: it does not need a ‘\0’ terminator.
long count
    maximum number of rows to return, or 0 for no limit

Return Value
See SPI_execute_plan.
SPI_processed and SPI_tuptable are set as in SPI_execute if successful.
SPI_cursor_open

Name
SPI_cursor_open — set up a cursor using a statement created with SPI_prepare

Synopsis
 Portal SPI_cursor_open(const char * name, SPIPlanPtr plan,
                   Datum * values, const char * nulls,
                   bool read_only)

Description
SPI_cursor_open sets up a cursor (internally, a portal) that will execute a statement prepared
by SPI_prepare. The parameters have the same meanings as the corresponding parameters to
SPI_execute_plan.

Using a cursor instead of executing the statement directly has two benefits. First, the result rows can
be retrieved a few at a time, avoiding memory overrun for queries that return many rows. Second,
a portal can outlive the current procedure (it can, in fact, live to the end of the current transaction).
Returning the portal name to the procedure’s caller provides a way of returning a row set as result.
The passed-in parameter data will be copied into the cursor’s portal, so it can be freed while the cursor
still exists.

Arguments
 const char * name
      name for portal, or NULL to let the system select a name
SPIPlanPtr plan
      prepared statement (returned by SPI_prepare)
 Datum * values
      An array of actual parameter values. Must have same length as the statement’s number of argu-
ments.
 const char * nulls
      An array describing which parameters are null. Must have same length as the statement’s number
of arguments.

If nulls is NULL then SPI_cursor_open assumes that no parameters are null. Otherwise, each entry of the nulls array should be ‘ ’ if the corresponding parameter value is non-null, or ‘n’ if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding values entry doesn’t matter.) Note that nulls is not a text string, just an array: it does not need a ‘\0’ terminator.
bool read_only
    true for read-only execution

Return Value

Pointer to portal containing the cursor. Note there is no error return convention; any error will be reported via elog.
SPI_cursor_open_with_args

Name

SPI_cursor_open_with_args — set up a cursor using a query and parameters

Synopsis

Portal SPI_cursor_open_with_args(const char * name, const char * command, int nargs, Oid * argtypes, Datum * values, const char * nulls, bool read_only, int cursorOptions)

Description

SPI_cursor_open_with_args sets up a cursor (internally, a portal) that will execute the specified query. Most of the parameters have the same meanings as the corresponding parameters to SPI_prepare_cursor and SPI_cursor_open.

For one-time query execution, this function should be preferred over SPI_prepare_cursor followed by SPI_cursor_open. If the same command is to be executed with many different parameters, either method might be faster, depending on the cost of re-planning versus the benefit of custom plans.

The passed-in parameter data will be copied into the cursor’s portal, so it can be freed while the cursor still exists.

Arguments

const char * name
    name for portal, or NULL to let the system select a name
const char * command
    command string
int nargs
    number of input parameters ($1, $2, etc.)
Oid * argtypes
    an array of length nargs, containing the OIDs of the data types of the parameters
Datum * values
    an array of length nargs, containing the actual parameter values
const char * nulls
    an array of length nargs, describing which parameters are null

If nulls is NULL then SPI_cursor_open_with_args assumes that no parameters are null. Otherwise, each entry of the nulls array should be ‘ ’ if the corresponding parameter value is
SPI_cursor_open_with_args

non-null, or ‘n’ if the corresponding parameter value is null. (In the latter case, the actual value in the corresponding values entry doesn’t matter.) Note that nulls is not a text string, just an array: it does not need a ‘\0’ terminator.

bool read_only
    true for read-only execution

int cursorOptions
    integer bit mask of cursor options; zero produces default behavior

Return Value

Pointer to portal containing the cursor. Note there is no error return convention; any error will be reported via elog.
**SPI_cursor_open_with_paramlist**

**Name**

SPI_cursor_open_with_paramlist — set up a cursor using parameters

**Synopsis**

Portal SPI_cursor_open_with_paramlist(const char *name,
    SPIPlanPtr plan,
    ParamListInfo params,
    bool read_only)

**Description**

SPI_cursor_open_with_paramlist sets up a cursor (internally, a portal) that will execute a statement prepared by SPI_prepare. This function is equivalent to SPI_cursor_open except that information about the parameter values to be passed to the query is presented differently. The ParamListInfo representation can be convenient for passing down values that are already available in that format. It also supports use of dynamic parameter sets via hook functions specified in ParamListInfo.

The passed-in parameter data will be copied into the cursor’s portal, so it can be freed while the cursor still exists.

**Arguments**

const char * name
    name for portal, or NULL to let the system select a name

SPIPlanPtr plan
    prepared statement (returned by SPI_prepare)

ParamListInfo params
    data structure containing parameter types and values; NULL if none

bool read_only
    true for read-only execution

**Return Value**

Pointer to portal containing the cursor. Note there is no error return convention; any error will be reported via elog.
SPI_cursor_find

Name
SPI_cursor_find — find an existing cursor by name

Synopsis
Portal SPI_cursor_find(const char * name)

Description
SPI_cursor_find finds an existing portal by name. This is primarily useful to resolve a cursor name returned as text by some other function.

Arguments
const char * name
   name of the portal

Return Value
pointer to the portal with the specified name, or NULL if none was found
SPI_cursor_fetch

Name
SPI_cursor_fetch — fetch some rows from a cursor

Synopsis
void SPI_cursor_fetch(Portal portal, bool forward, long count)

Description
SPI_cursor_fetch fetches some rows from a cursor. This is equivalent to a subset of the SQL command FETCH (see SPI_scroll_cursor_fetch for more functionality).

Arguments
Portal portal
    portal containing the cursor
bool forward
    true for fetch forward, false for fetch backward
long count
    maximum number of rows to fetch

Return Value
SPI_processed and SPI_tuptable are set as in SPI_execute if successful.

Notes
Fetching backward may fail if the cursor’s plan was not created with the CURSOR_OPT_SCROLL option.
SPI_cursor_move

Name

SPI_cursor_move — move a cursor

Synopsis

void SPI_cursor_move(Portal portal, bool forward, long count)

Description

SPI_cursor_move skips over some number of rows in a cursor. This is equivalent to a subset of the SQL command MOVE (see SPI_scroll_cursor_move for more functionality).

Arguments

Portal portal
    portal containing the cursor
bool forward
    true for move forward, false for move backward
long count
    maximum number of rows to move

Notes

Moving backward may fail if the cursor’s plan was not created with the CURSOR_OPT_SCROLL option.
SPI_scroll_cursor_fetch

**Name**

SPI_scroll_cursor_fetch — fetch some rows from a cursor

**Synopsis**

```c
void SPI_scroll_cursor_fetch(Portal portal, FetchDirection direction,
                            long count)
```

**Description**

SPI_scroll_cursor_fetch fetches some rows from a cursor. This is equivalent to the SQL command `FETCH`.

**Arguments**

- **Portal portal**
  portal containing the cursor
- **FetchDirection direction**
  one of `FETCH_FORWARD`, `FETCH_BACKWARD`, `FETCH_ABSOLUTE` or `FETCH_RELATIVE`
- **long count**
  number of rows to fetch for `FETCH_FORWARD` or `FETCH_BACKWARD`; absolute row number to fetch for `FETCH_ABSOLUTE`; or relative row number to fetch for `FETCH_RELATIVE`

**Return Value**

SPI_processed and SPI_tuptable are set as in SPI_execute if successful.

**Notes**

See the SQL FETCH command for details of the interpretation of the `direction` and `count` parameters.

Direction values other than `FETCH_FORWARD` may fail if the cursor’s plan was not created with the CURSOR_OPT_SCROLL option.
SPI_scroll_cursor_move

Name
SPI_scroll_cursor_move — move a cursor

Synopsis

void SPI_scroll_cursor_move(Portal portal, FetchDirection direction,
   long count)

Description

SPI_scroll_cursor_move skips over some number of rows in a cursor. This is equivalent to the SQL command MOVE.

Arguments

Portal portal
   portal containing the cursor
FetchDirection direction
   one of FETCH_FORWARD, FETCH_BACKWARD, FETCH_ABSOLUTE or FETCH_RELATIVE
long count
   number of rows to move for FETCH_FORWARD or FETCH_BACKWARD; absolute row number to move to for FETCH_ABSOLUTE; or relative row number to move to for FETCH_RELATIVE

Return Value

SPI_processed is set as in SPI_execute if successful. SPI_tuptable is set to NULL, since no rows are returned by this function.

Notes

See the SQL FETCH command for details of the interpretation of the direction and count parameters.

Direction values other than FETCH_FORWARD may fail if the cursor’s plan was not created with the CURSOR_OPT_SCROLL option.
SPI_cursor_close

Name
SPI_cursor_close — close a cursor

Synopsis
void SPI_cursor_close(Portal portal)

Description
SPI_cursor_close closes a previously created cursor and releases its portal storage.
All open cursors are closed automatically at the end of a transaction. SPI_cursor_close need only
be invoked if it is desirable to release resources sooner.

Arguments
Portal portal
    portal containing the cursor
SPI_keepplan

Name
SPI_keepplan — save a prepared statement

Synopsis
int SPI_keepplan(SPIPlanPtr plan)

Description
SPI_keepplan saves a passed statement (prepared by SPI_prepare) so that it will not be freed by
SPI_finish nor by the transaction manager. This gives you the ability to reuse prepared statements
in the subsequent invocations of your procedure in the current session.

Arguments
SPIPlanPtr plan
    the prepared statement to be saved

Return Value
0 on success; SPI_ERROR_ARGUMENT if plan is NULL or invalid

Notes
The passed-in statement is relocated to permanent storage by means of pointer adjustment (no data
copying is required). If you later wish to delete it, use SPI_freeplan on it.
**SPI_saveplan**

**Name**

SPI_saveplan — save a prepared statement

**Synopsis**

SPIPlanPtr SPI_saveplan(SPIPlanPtr plan)

**Description**

SPI_saveplan copies a passed statement (prepared by SPI_prepare) into memory that will not be freed by SPI_finish nor by the transaction manager, and returns a pointer to the copied statement. This gives you the ability to reuse prepared statements in the subsequent invocations of your procedure in the current session.

**Arguments**

SPIPlanPtr plan

the prepared statement to be saved

**Return Value**

Pointer to the copied statement; or NULL if unsuccessful. On error, SPI_result is set thus:

- **SPI_ERROR_ARGUMENT**
  - if plan is NULL or invalid
- **SPI_ERROR_UNCONNECTED**
  - if called from an unconnected procedure

**Notes**

The originally passed-in statement is not freed, so you might wish to do SPI_freeplan on it to avoid leaking memory until SPI_finish.

In most cases, SPI_keepplan is preferred to this function, since it accomplishes largely the same result without needing to physically copy the prepared statement’s data structures.
45.2. Interface Support Functions

The functions described here provide an interface for extracting information from result sets returned by SPI_execute and other SPI functions.

All functions described in this section can be used by both connected and unconnected procedures.

**SPI_fname**

**Name**

SPI_fname — determine the column name for the specified column number

**Synopsis**

```c
char * SPI_fname(TupleDesc rowdesc, int colnumber)
```

**Description**

SPI_fname returns a copy of the column name of the specified column. (You can use pfree to release the copy of the name when you don’t need it anymore.)

**Arguments**

- TupleDesc rowdesc
  - input row description
- int colnumber
  - column number (count starts at 1)

**Return Value**

The column name; NULL if colnumber is out of range. SPI_result set to SPI_ERROR_NOATTRIBUTE on error.
SPI_fnumber

Name
SPI_fnumber — determine the column number for the specified column name

Synopsis
int SPI_fnumber(TupleDesc rowdesc, const char * colname)

Description
SPI_fnumber returns the column number for the column with the specified name.

If colname refers to a system column (e.g., oid) then the appropriate negative column number
will be returned. The caller should be careful to test the return value for exact equality to
SPI_ERROR_NOATTRIBUTE to detect an error; testing the result for less than or equal to 0 is not
correct unless system columns should be rejected.

Arguments

TupleDesc rowdesc
input row description

const char * colname
column name

Return Value
Column number (count starts at 1), or SPI_ERROR_NOATTRIBUTE if the named column was not
found.
**SPI_getvalue**

**Name**

SPI_getvalue — return the string value of the specified column

**Synopsis**

```c
char * SPI_getvalue(HeapTuple row, TupleDesc rowdesc, int colnumber)
```

**Description**

SPI_getvalue returns the string representation of the value of the specified column. The result is returned in memory allocated using `palloc`. (You can use `pfree` to release the memory when you don’t need it anymore.)

**Arguments**

- `HeapTuple row` 
  input row to be examined
- `TupleDesc rowdesc` 
  input row description
- `int colnumber` 
  column number (count starts at 1)

**Return Value**

Column value, or `NULL` if the column is null, `colnumber` is out of range (SPI_result is set to SPI_ERROR_NOATTRIBUTE), or no output function is available (SPI_result is set to SPI_ERROR_NOOUTFUNC).
SPI_getbinval

Name
SPI_getbinval — return the binary value of the specified column

Synopsis
Datum SPI_getbinval(HeapTuple row, TupleDesc rowdesc, int colnumber,
     bool * isnull)

Description
SPI_getbinval returns the value of the specified column in the internal form (as type Datum).
This function does not allocate new space for the datum. In the case of a pass-by-reference data type,
the return value will be a pointer into the passed row.

Arguments
HeapTuple row
    input row to be examined
TupleDesc rowdesc
    input row description
int colnumber
    column number (count starts at 1)
bool * isnull
    flag for a null value in the column

Return Value
The binary value of the column is returned. The variable pointed to by isnull is set to true if the
column is null, else to false.

SPI_result is set to SPI_ERROR_NOATTRIBUTE on error.
SPI_gettype

Name
SPI_gettype — return the data type name of the specified column

Synopsis
char * SPI_gettype(TupleDesc rowdesc, int colnumber)

Description
SPI_gettype returns a copy of the data type name of the specified column. (You can use pfree to release the copy of the name when you don’t need it anymore.)

Arguments

TupleDesc rowdesc
   input row description

int colnumber
   column number (count starts at 1)

Return Value
The data type name of the specified column, or NULL on error. SPI_result is set to SPI_ERROR_NOATTRIBUTE on error.
SPI_gettypeid

Name
SPI_gettypeid — return the data type OID of the specified column

Synopsis
Oid SPI_gettypeid(TupleDesc rowdesc, int colnumber)

Description
SPI_gettypeid returns the OID of the data type of the specified column.

Arguments

TupleDesc rowdesc
    input row description
int colnumber
    column number (count starts at 1)

Return Value
The OID of the data type of the specified column or InvalidOid on error. On error, SPI_result is set to SPI_ERROR_NOATTRIBUTE.
**SPI_getrelname**

**Name**
SPI_getrelname — return the name of the specified relation

**Synopsis**
```c
char * SPI_getrelname(Relation rel)
```

**Description**
SPI_getrelname returns a copy of the name of the specified relation. (You can use pfree to release the copy of the name when you don’t need it anymore.)

**Arguments**
Relation rel
   input relation

**Return Value**
The name of the specified relation.
**SPI_getnspname**

**Name**

SPI_getnspname — return the namespace of the specified relation

**Synopsis**

char * SPI_getnspname(Relation rel)

**Description**

SPI_getnspname returns a copy of the name of the namespace that the specified Relation belongs to. This is equivalent to the relation’s schema. You should `pfree` the return value of this function when you are finished with it.

**Arguments**

Relation rel

    input relation

**Return Value**

The name of the specified relation’s namespace.
45.3. Memory Management

PostgreSQL allocates memory within memory contexts, which provide a convenient method of managing allocations made in many different places that need to live for differing amounts of time. Destroying a context releases all the memory that was allocated in it. Thus, it is not necessary to keep track of individual objects to avoid memory leaks; instead only a relatively small number of contexts have to be managed. `palloc` and related functions allocate memory from the “current” context.

`SPI_connect` creates a new memory context and makes it current. `SPI_finish` restores the previous current memory context and destroys the context created by `SPI_connect`. These actions ensure that transient memory allocations made inside your procedure are reclaimed at procedure exit, avoiding memory leakage.

However, if your procedure needs to return an object in allocated memory (such as a value of a pass-by-reference data type), you cannot allocate that memory using `palloc`, at least not while you are connected to SPI. If you try, the object will be deallocated by `SPI_finish`, and your procedure will not work reliably. To solve this problem, use `SPI_palloc` to allocate memory for your return object. `SPI_palloc` allocates memory in the “upper executor context”, that is, the memory context that was current when `SPI_connect` was called, which is precisely the right context for a value returned from your procedure.

If `SPI_palloc` is called while the procedure is not connected to SPI, then it acts the same as a normal `palloc`. Before a procedure connects to the SPI manager, the current memory context is the upper executor context, so all allocations made by the procedure via `palloc` or by SPI utility functions are made in this context.

When `SPI_connect` is called, the private context of the procedure, which is created by `SPI_connect`, is made the current context. All allocations made by `palloc`, `repalloc`, or SPI utility functions (except for `SPI_copytuple`, `SPI_returntuple`, `SPI_modifytuple`, and `SPI_palloc`) are made in this context. When a procedure disconnects from the SPI manager (via `SPI_finish`) the current context is restored to the upper executor context, and all allocations made in the procedure memory context are freed and cannot be used any more.

All functions described in this section can be used by both connected and unconnected procedures. In an unconnected procedure, they act the same as the underlying ordinary server functions (`palloc`, etc.).

### SPI_palloc

**Name**

SPI_palloc — allocate memory in the upper executor context

**Synopsis**

```c
void * SPI_palloc(Size size)
```

**Description**

`SPI_palloc` allocates memory in the upper executor context.
Arguments

Size size

size in bytes of storage to allocate

Return Value

pointer to new storage space of the specified size
SPI_repalloclip

Name
SPI_repalloclip — reallocate memory in the upper executor context

Synopsis
void * SPI_repalloclip(void * pointer, Size size)

Description
SPI_repalloclip changes the size of a memory segment previously allocated using SPI_palloc. This function is no longer different from plain repalloc. It’s kept just for backward compatibility of existing code.

Arguments
void * pointer
    pointer to existing storage to change
Size size
    size in bytes of storage to allocate

Return Value
pointer to new storage space of specified size with the contents copied from the existing area
SPI_pfree

Name
SPI_pfree — free memory in the upper executor context

Synopsis
void SPI_pfree(void * pointer)

Description
SPI_pfree frees memory previously allocated using SPI_palloc or SPI_repalloc.
This function is no longer different from plain pfree. It’s kept just for backward compatibility of existing code.

Arguments
void * pointer
    pointer to existing storage to free
SPI_copytuple

Name
SPI_copytuple — make a copy of a row in the upper executor context

Synopsis
HeapTuple SPI_copytuple(HeapTuple row)

Description
SPI_copytuple makes a copy of a row in the upper executor context. This is normally used to return a modified row from a trigger. In a function declared to return a composite type, use SPI_returntuple instead.

Arguments
HeapTuple row
row to be copied

Return Value
the copied row; NULL only if tuple is NULL
**SPI_returntuple**

**Name**

SPI_returntuple — prepare to return a tuple as a Datum

**Synopsis**

HeapTupleHeader SPI_returntuple(HeapTuple row, TupleDesc rowdesc)

**Description**

SPI_returntuple makes a copy of a row in the upper executor context, returning it in the form of a row type Datum. The returned pointer need only be converted to Datum via PointerGetDatum before returning.

Note that this should be used for functions that are declared to return composite types. It is not used for triggers; use SPI_copytuple for returning a modified row in a trigger.

**Arguments**

HeapTuple row
- row to be copied

TupleDesc rowdesc
- descriptor for row (pass the same descriptor each time for most effective caching)

**Return Value**

HeapTupleHeader pointing to copied row; NULL only if row or rowdesc is NULL
**SPI_modifytuple**

**Name**

SPI_modifytuple — create a row by replacing selected fields of a given row

**Synopsis**

HeapTuple SPI_modifytuple(Relation rel, HeapTuple row, int ncols, int * colnum, Datum * values, const char * nulls)

**Description**

SPI_modifytuple creates a new row by substituting new values for selected columns, copying the original row’s columns at other positions. The input row is not modified.

**Arguments**

Relation rel

Used only as the source of the row descriptor for the row. (Passing a relation rather than a row descriptor is a misfeature.)

HeapTuple row

row to be modified

int ncols

number of columns to be changed

int * colnum

an array of length ncols, containing the numbers of the columns that are to be changed (column numbers start at 1)

Datum * values

an array of length ncols, containing the new values for the specified columns

const char * nulls

an array of length ncols, describing which new values are null

If nulls is NULL then SPI_modifytuple assumes that no new values are null. Otherwise, each entry of the nulls array should be ‘ ’ if the corresponding new value is non-null, or ‘n’ if the corresponding new value is null. (In the latter case, the actual value in the corresponding values entry doesn’t matter.) Note that nulls is not a text string, just an array: it does not need a ‘\0’ terminator.

**Return Value**

new row with modifications, allocated in the upper executor context; NULL only if row is NULL
On error, SPI_result is set as follows:

**SPI_ERROR_ARGUMENT**
- if rel is NULL, or if row is NULL, or if ncols is less than or equal to 0, or if colnum is NULL, or if values is NULL.

**SPI_ERROR_NOATTRIBUTE**
- if colnum contains an invalid column number (less than or equal to 0 or greater than the number of column in row)
SPI_freetuple

Name
SPI_freetuple — free a row allocated in the upper executor context

Synopsis
void SPI_freetuple(HeapTuple row)

Description
SPI_freetuple frees a row previously allocated in the upper executor context.
This function is no longer different from plain heap_freetuple. It’s kept just for backward compatibility of existing code.

Arguments
HeapTuple row
row to free
SPI_freetuptable

Name
SPI_freetuptable — free a row set created by SPI_execute or a similar function

Synopsis
void SPI_freetuptable(SPI_tupleTable * tuptable)

Description
SPI_freetuptable frees a row set created by a prior SPI command execution function, such as SPI_execute. Therefore, this function is often called with the global variable SPI_tuptable as argument.

This function is useful if a SPI procedure needs to execute multiple commands and does not want to keep the results of earlier commands around until it ends. Note that any unfreed row sets will be freed anyway at SPI_finish. Also, if a subtransaction is started and then aborted within execution of a SPI procedure, SPI automatically frees any row sets created while the subtransaction was running.

Beginning in PostgreSQL 9.3, SPI_freetuptable contains guard logic to protect against duplicate deletion requests for the same row set. In previous releases, duplicate deletions would lead to crashes.

Arguments
SPI_tupleTable * tuptable
    pointer to row set to free, or NULL to do nothing
SPI_freeplan

Name
SPI_freeplan — free a previously saved prepared statement

Synopsis
int SPI_freeplan(SPIPlanPtr plan)

Description
SPI_freeplan releases a prepared statement previously returned by SPI_prepare or saved by SPI_keepplan or SPI_saveplan.

Arguments
SPIPlanPtr plan
        pointer to statement to free

Return Value
0 on success; SPI_ERROR_ARGUMENT if plan is NULL or invalid
45.4. Visibility of Data Changes

The following rules govern the visibility of data changes in functions that use SPI (or any other C function):

- During the execution of an SQL command, any data changes made by the command are invisible to the command itself. For example, in:
  
  ```sql
  INSERT INTO a SELECT * FROM a;
  ```
  the inserted rows are invisible to the `SELECT` part.

- Changes made by a command C are visible to all commands that are started after C, no matter whether they are started inside C (during the execution of C) or after C is done.

- Commands executed via SPI inside a function called by an SQL command (either an ordinary function or a trigger) follow one or the other of the above rules depending on the read/write flag passed to SPI. Commands executed in read-only mode follow the first rule: they cannot see changes of the calling command. Commands executed in read-write mode follow the second rule: they can see all changes made so far.

- All standard procedural languages set the SPI read-write mode depending on the volatility attribute of the function. Commands of `STABLE` and `IMMUTABLE` functions are done in read-only mode, while commands of `VOLATILE` functions are done in read-write mode. While authors of C functions are able to violate this convention, it’s unlikely to be a good idea to do so.

The next section contains an example that illustrates the application of these rules.

45.5. Examples

This section contains a very simple example of SPI usage. The procedure `execq` takes an SQL command as its first argument and a row count as its second, executes the command using `SPI_exec` and returns the number of rows that were processed by the command. You can find more complex examples for SPI in the source tree in `src/test/regress/regress.c` and in the `spi` module.

```c
#include "postgres.h"
#include "executor/spi.h"
#include "utils/builtins.h"

#ifdef PG_MODULE_MAGIC
PG_MODULE_MAGIC;
#endif

int64 execq(text *sql, int cnt);

int64
execq(text *sql, int cnt)
{
  char *command;
  int ret;
  uint64 proc;

  /* Convert given text object to a C string */
```
command = text_to_cstring(sql);

SPI_connect();

ret = SPI_exec(command, cnt);

proc = SPI_processed;
/*
 * If some rows were fetched, print them via elog(INFO).
 */
if (ret > 0 && SPI_tuptable != NULL)
{
    TupleDesc tupdesc = SPI_tuptable->tupdesc;
    SPITupleTable *tuptable = SPI_tuptable;
    char buf[8192];
    uint64 j;

    for (j = 0; j < proc; j++)
    {
        HeapTuple tuple = tuptable->vals[j];
        int i;

        for (i = 1, buf[0] = 0; i <= tupdesc->natts; i++)
            snprintf(buf + strlen(buf), sizeof(buf) - strlen(buf), " %s%s",
                SPI_getvalue(tuple, tupdesc, i),
                (i == tupdesc->natts) ? " " : " |");
        elog(INFO, "EXECQ: %s", buf);
    }
}

SPI_finish();
pfree(command);

return (proc);
}

(This function uses call convention version 0, to make the example easier to understand. In real applications you should use the new version 1 interface.)

This is how you declare the function after having compiled it into a shared library (details are in Section 36.9.6.):

CREATE FUNCTION execq(text, integer) RETURNS int8
    AS 'filename'
    LANGUAGE C STRICT;

Here is a sample session:

=> SELECT execq('CREATE TABLE a (x integer)', 0);
execq
-------
 0
(1 row)

=> INSERT INTO a VALUES (execq('INSERT INTO a VALUES (0)', 0));
INSERT 0 1
-> SELECT execq('SELECT * FROM a', 0);
INFO: EXECQ: 0 -- inserted by execq
INFO: EXECQ: 1 -- returned by execq and inserted by upper INSERT
execq
-------
2
(1 row)

-> SELECT execq('INSERT INTO a SELECT x + 2 FROM a', 1);
execq
-------
1
(1 row)

-> SELECT execq('SELECT * FROM a', 10);
INFO: EXECQ: 0
INFO: EXECQ: 1
INFO: EXECQ: 2 -- 0 + 2, only one row inserted - as specified
execq
-------
3 -- 10 is the max value only, 3 is the real number of rows
(1 row)

-> DELETE FROM a;
DELETE 3
-> INSERT INTO a VALUES (execq('SELECT * FROM a', 0) + 1);
INSERT 0 1
-> SELECT * FROM a;
x
---
1 -- no rows in a (0) + 1
(1 row)

-> INSERT INTO a VALUES (execq('SELECT * FROM a', 0) + 1);
INFO: EXECQ: 1
INSERT 0 1
-> SELECT * FROM a;
x
---
1
2 -- there was one row in a + 1
(2 rows)

-- This demonstrates the data changes visibility rule:

-> INSERT INTO a SELECT execq('SELECT * FROM a', 0) * x FROM a;
INFO: EXECQ: 1
INFO: EXECQ: 2
INFO: EXECQ: 1
INFO: EXECQ: 2
INFO: EXECQ: 2
INSERT 0 2
-> SELECT * FROM a;
x
---
Chapter 45. Server Programming Interface

1

2

-- 2 rows * 1 (x in first row)

6

-- 3 rows (2 + 1 just inserted) * 2 (x in second row)

(4 rows)

^^^^^^

rows visible to execq() in different invocations
Chapter 46. Background Worker Processes

PostgreSQL can be extended to run user-supplied code in separate processes. Such processes are started, stopped and monitored by `postgres`, which permits them to have a lifetime closely linked to the server’s status. These processes have the option to attach to PostgreSQL’s shared memory area and to connect to databases internally; they can also run multiple transactions serially, just like a regular client-connected server process. Also, by linking to libpq they can connect to the server and behave like a regular client application.

Warning

There are considerable robustness and security risks in using background worker processes because, being written in the C language, they have unrestricted access to data. Administrators wishing to enable modules that include background worker processes should exercise extreme caution. Only carefully audited modules should be permitted to run background worker processes.

Background workers can be initialized at the time that PostgreSQL is started by including the module name in `shared_preload_libraries`. A module wishing to run a background worker can register it by calling `RegisterBackgroundWorker(BackgroundWorker *worker)` from its `_PG_init()`. Background workers can also be started after the system is up and running by calling the function `RegisterDynamicBackgroundWorker(BackgroundWorker *worker, BackgroundWorkerHandle **handle)`. Unlike `RegisterBackgroundWorker`, which can only be called from within the postmaster, `RegisterDynamicBackgroundWorker` must be called from a regular backend.

The structure `BackgroundWorker` is defined thus:

```c
typedef void (*bgworker_main_type)(Datum main_arg);
typedef struct BackgroundWorker
{
    char bgw_name[BGW_MAXLEN];
    int bgw_flags;
    BgWorkerStartTime bgw_start_time;
    int bgw_restart_time; /* in seconds, or BGW_NEVER_RESTART */
    bgworker_main_type bgw_main;
    char bgw_library_name[BGW_MAXLEN]; /* only if bgw_main is NULL */
    char bgw_function_name[BGW_MAXLEN]; /* only if bgw_main is NULL */
    Datum bgw_main_arg;
    char bgw_extra[BGW_EXTRALEN];
    int bgw_notify_pid;
} BackgroundWorker;
```

`bgw_name` is a string to be used in log messages, process listings and similar contexts.

`bgw_flags` is a bitwise-or’d bit mask indicating the capabilities that the module wants. Possible values are:

**BGWORKER_SHMEM_ACCESS**

Requests shared memory access. Workers without shared memory access cannot access any of PostgreSQL’s shared data structures, such as heavyweight or lightweight locks, shared buffers, or any custom data structures which the worker itself may wish to create and use.
Chapter 46. Background Worker Processes

BGWORKER_BACKEND_DATABASE_CONNECTION

Requests the ability to establish a database connection through which it can later run transactions and queries. A background worker using BGWORKER_BACKEND_DATABASE_CONNECTION to connect to a database must also attach shared memory using BGWORKER_SHMEM_ACCESS, or worker start-up will fail.

bgw_start_time is the server state during which postgres should start the process; it can be one of BgWorkerStart_PostmasterStart (start as soon as postgres itself has finished its own initialization; processes requesting this are not eligible for database connections), BgWorkerStart_ConsistentState (start as soon as a consistent state has been reached in a hot standby, allowing processes to connect to databases and run read-only queries), and BgWorkerStart_RecoveryFinished (start as soon as the system has entered normal read-write state). Note the last two values are equivalent in a server that’s not a hot standby. Note that this setting only indicates when the processes are to be started; they do not stop when a different state is reached.

bgw_restart_time is the interval, in seconds, that postgres should wait before restarting the process, in case it crashes. It can be any positive value, or BGW_NEVER_RESTART, indicating not to restart the process in case of a crash.

bgw_main is a pointer to the function to run when the process is started. This field can only safely be used to launch functions within the core server, because shared libraries may be loaded at different starting addresses in different backend processes. This will happen on all platforms when the library is loaded using any mechanism other than shared_preload_libraries. Even when that mechanism is used, address space layout variations will still occur on Windows, and when EXEC_BACKEND is used. Therefore, most users of this API should set this field to NULL. If it is non-NULL, it takes precedence over bgw_library_name and bgw_function_name.

bgw_library_name is the name of a library in which the initial entry point for the background worker should be sought. The named library will be dynamically loaded by the worker process and bgw_function_name will be used to identify the function to be called. If loading a function from the core code, bgw_main should be set instead.

bgw_function_name is the name of a function in a dynamically loaded library which should be used as the initial entry point for a new background worker.

bgw_main_arg is the Datum argument to the background worker main function. Regardless of whether that function is specified via bgw_main or via the combination of bgw_library_name and bgw_function_name, this main function should take a single argument of type Datum and return void. bgw_main_arg will be passed as the argument. In addition, the global variable MyBgworkerEntry points to a copy of the BackgroundWorker structure passed at registration time; the worker may find it helpful to examine this structure.

On Windows (and anywhere else where EXEC_BACKEND is defined) or in dynamic background workers it is not safe to pass a Datum by reference, only by value. If an argument is required, it is safest to pass an int32 or other small value and use that as an index into an array allocated in shared memory. If a value like a cstring or text is passed then the pointer won’t be valid from the new background worker process.

bgw_extra can contain extra data to be passed to the background worker. Unlike bgw_main_arg, this data is not passed as an argument to the worker’s main function, but it can be accessed via MyBgworkerEntry, as discussed above.

bgw_notify_pid is the PID of a PostgreSQL backend process to which the postmaster should send SIGUSR1 when the process is started or exits. It should be 0 for workers registered at postmaster
startup time, or when the backend registering the worker does not wish to wait for the worker to start up. Otherwise, it should be initialized to MyProcPid.

Once running, the process can connect to a database by calling 
\texttt{BackgroundWorkerInitializeConnection(char *dbname, char *username)} or 
\texttt{BackgroundWorkerInitializeConnectionByOid(Oid dboid, Oid useroid)}. This allows the process to run transactions and queries using the SPI interface. If \texttt{dbname} is \texttt{NULL} or \texttt{dboid} is \texttt{InvalidOid}, the session is not connected to any particular database, but shared catalogs can be accessed. If \texttt{username} is \texttt{NULL} or \texttt{useroid} is \texttt{InvalidOid}, the process will run as the superuser created during \texttt{initdb}. A background worker can only call one of these two functions, and only once. It is not possible to switch databases.

Signals are initially blocked when control reaches the \texttt{bgw_main} function, and must be unblocked by it; this is to allow the process to customize its signal handlers, if necessary. Signals can be unblocked in the new process by calling \texttt{BackgroundWorkerUnblockSignals} and blocked by calling \texttt{BackgroundWorkerBlockSignals}.

If \texttt{bgw_restart_time} for a background worker is configured as \texttt{BGW_NEVER_RESTART}, or if it exits with an exit code of 0 or is terminated by \texttt{TerminateBackgroundWorker}, it will be automatically unregistered by the postmaster on exit. Otherwise, it will be restarted after the time period configured via \texttt{bgw_restart_time}, or immediately if the postmaster reinitializes the cluster due to a backend failure. Backends which need to suspend execution only temporarily should use an interruptible sleep rather than exiting; this can be achieved by calling \texttt{WaitLatch()}. Make sure the \texttt{WL_POSTMASTER_DEATH} flag is set when calling that function, and verify the return code for a prompt exit in the emergency case that \texttt{postgres} itself has terminated.

When a background worker is registered using the \texttt{RegisterDynamicBackgroundWorker} function, it is possible for the backend performing the registration to obtain information regarding the status of the worker. Backends wishing to do this should pass the address of a \texttt{BackgroundWorkerHandle *} as the second argument to \texttt{RegisterDynamicBackgroundWorker}. If the worker is successfully registered, this pointer will be initialized with an opaque handle that can subsequently be passed to \texttt{GetBackgroundWorkerPid(BackgroundWorkerHandle *, pid_t *)} or \texttt{TerminateBackgroundWorker(BackgroundWorkerHandle *)}. \texttt{GetBackgroundWorkerPid} can be used to poll the status of the worker; a return value of \texttt{BGWH_NOT_YET_STARTED} indicates that the worker has not yet been started by the postmaster; \texttt{BGWH_STOPPED} indicates that it has been started but is no longer running; and \texttt{BGWH_STARTED} indicates that it is currently running. In this last case, the PID will also be returned via the second argument. \texttt{TerminateBackgroundWorker} causes the postmaster to send \texttt{SIGTERM} to the worker if it is running, and to unregister it as soon as it is not.

In some cases, a process which registers a background worker may wish to wait for the worker to start up. This can be accomplished by initializing \texttt{bgw_notify_pid} to \texttt{MyProcPid} and then passing the \texttt{BackgroundWorkerHandle *} obtained at registration time to \texttt{WaitForBackgroundWorkerStartup(BackgroundWorkerHandle *handle, pid_t *)} function. This function will block until the postmaster has attempted to start the background worker, or until the postmaster dies. If the background runner is running, the return value will be \texttt{BGWH_STARTED}, and the PID will be written to the provided address. Otherwise, the return value will be \texttt{BGWH_STOPPED} or \texttt{BGWH_POSTMASTER_DIED}.

If a background worker sends asynchronous notifications with the \texttt{NOTIFY} command via the Server Programming Interface (SPI), it should call \texttt{ProcessCompletedNotifies} explicitly after committing the enclosing transaction so that any notifications can be delivered. If a background worker registers to receive asynchronous notifications with the \texttt{LISTEN} through SPI, the worker will log those notifications, but there is no programmatic way for the worker to intercept and respond to those notifications.
The `src/test/modules/worker_spi` module contains a working example, which demonstrates some useful techniques.

The maximum number of registered background workers is limited by `max_worker_processes`. 
Chapter 47. Logical Decoding

PostgreSQL provides infrastructure to stream the modifications performed via SQL to external consumers. This functionality can be used for a variety of purposes, including replication solutions and auditing.

Changes are sent out in streams identified by logical replication slots.

The format in which those changes are streamed is determined by the output plugin used. An example plugin is provided in the PostgreSQL distribution. Additional plugins can be written to extend the choice of available formats without modifying any core code. Every output plugin has access to each individual new row produced by \texttt{INSERT} and the new row version created by \texttt{UPDATE}. Availability of old row versions for \texttt{UPDATE} and \texttt{DELETE} depends on the configured replica identity (see \texttt{REPLICA IDENTITY}).

Changes can be consumed either using the streaming replication protocol (see Section 51.3 and Section 47.3), or by calling functions via SQL (see Section 47.4). It is also possible to write additional methods of consuming the output of a replication slot without modifying core code (see Section 47.7).

47.1. Logical Decoding Examples

The following example demonstrates controlling logical decoding using the SQL interface.

Before you can use logical decoding, you must set \texttt{wal_level} to \texttt{logical} and \texttt{max_replication_slots} to at least 1. Then, you should connect to the target database (in the example below, \texttt{postgres}) as a superuser.

```sql
postgres=# -- Create a slot named 'regression_slot' using the output plugin 'test_decoding'
postgres=# SELECT * FROM pg_create_logical_replication_slot('regression_slot', 'test_decoding');

slot_name | xlog_position
-----------------+---------------
regression_slot | 0/16B1970
(1 row)

postgres=# SELECT slot_name, plugin, slot_type, database, active, restart_lsn, confirmed_flush_lsn FROM pg_replication_slots;

slot_name | plugin | slot_type | database | active | restart_lsn | confirmed_flush_lsn
-----------------+---------------+-----------+----------+--------+-------------+-----------------
regression_slot | test_decoding | logical | postgres | f | 0/16A4408 | 0/16A4440
(1 row)

postgres=# -- There are no changes to see yet
postgres=# SELECT * FROM pg_logical_slot_get_changes('regression_slot', NULL, NULL);

location | xid | data
----------+-----+------
(0 rows)

postgres=# CREATE TABLE data(id serial primary key, data text);

CREATE TABLE

postgres=# -- DDL isn’t replicated, so all you’ll see is the transaction
postgres=# SELECT * FROM pg_logical_slot_get_changes('regression_slot', NULL, NULL);

location | xid | data
----------+-----+------
0/16D5D48 | 688 | BEGIN 688
0/16E0380 | 688 | COMMIT 688
```

1282
Chapter 47. Logical Decoding

postgres=# -- Once changes are read, they're consumed and not emitted
postgres=# -- in a subsequent call:
postgres=# SELECT * FROM pg_logical_slot_get_changes('regression_slot', NULL, NULL);
location | xid | data
----------+-----+------
(0 rows)

postgres=# BEGIN;
postgres=# INSERT INTO data(data) VALUES('1');
postgres=# INSERT INTO data(data) VALUES('2');
postgres=# COMMIT;

postgres=# SELECT * FROM pg_logical_slot_get_changes('regression_slot', NULL, NULL);
location | xid | data
-----------+-----+-----------------------------------------------
0/16E0478 | 689 | BEGIN 689
0/16E0478 | 689 | table public.data: INSERT: id[integer]:1 data[text]:'1'
0/16E0580 | 689 | table public.data: INSERT: id[integer]:2 data[text]:'2'
0/16E0650 | 689 | COMMIT 689
(4 rows)

postgres=# INSERT INTO data(data) VALUES('3');

postgres=# -- You can also peek ahead in the change stream without consuming changes
postgres=# SELECT * FROM pg_logical_slot_peek_changes('regression_slot', NULL, NULL);
location | xid | data
-----------+-----+-----------------------------------------------
0/16E09C0 | 690 | BEGIN 690
0/16E09C0 | 690 | table public.data: INSERT: id[integer]:3 data[text]:'3'
0/16E0B90 | 690 | COMMIT 690
(3 rows)

postgres=# -- The next call to pg_logical_slot_peek_changes() returns the same changes as the previous one.
postgres=# SELECT * FROM pg_logical_slot_peek_changes('regression_slot', NULL, NULL);
location | xid | data
-----------+-----+-----------------------------------------------
0/16E09C0 | 690 | BEGIN 690
0/16E09C0 | 690 | table public.data: INSERT: id[integer]:3 data[text]:'3'
0/16E0B90 | 690 | COMMIT 690
(3 rows)

postgres=# -- options can be passed to output plugin, to influence the formatting
postgres=# SELECT * FROM pg_logical_slot_peek_changes('regression_slot', NULL, NULL, 'include-timestamp', 'on');
location | xid | data
-----------+-----+-----------------------------------------------
0/16E09C0 | 690 | BEGIN 690
0/16E09C0 | 690 | table public.data: INSERT: id[integer]:3 data[text]:'3'
0/16E0B90 | 690 | COMMIT 690 (at 2014-02-27 16:41:51.863092+01)
(3 rows)

postgres=# -- Remember to destroy a slot you no longer need to stop it consuming server resources:
postgres=# SELECT pg_drop_replication_slot('regression_slot');
Chapter 47. Logical Decoding

The following example shows how logical decoding is controlled over the streaming replication protocol, using the program pg_recvlogical included in the PostgreSQL distribution. This requires that client authentication is set up to allow replication connections (see Section 26.2.5.1) and that max_wal_senders is set sufficiently high to allow an additional connection.

```
$ pg_recvlogical -d postgres --slot test --create-slot
$ pg_recvlogical -d postgres --slot test --start -f -
$ Control+Z
$ psql -d postgres -c "INSERT INTO data(data) VALUES('4');"
$ fg
BEGIN 693
  table public.data: INSERT: id[integer]:4 data[text]:'4'
COMMIT 693
$ Control+C
$ pg_recvlogical -d postgres --slot test --drop-slot
```

47.2. Logical Decoding Concepts

47.2.1. Logical Decoding

Logical decoding is the process of extracting all persistent changes to a database’s tables into a coherent, easy to understand format which can be interpreted without detailed knowledge of the database’s internal state.

In PostgreSQL, logical decoding is implemented by decoding the contents of the write-ahead log, which describe changes on a storage level, into an application-specific form such as a stream of tuples or SQL statements.

47.2.2. Replication Slots

In the context of logical replication, a slot represents a stream of changes that can be replayed to a client in the order they were made on the origin server. Each slot streams a sequence of changes from a single database.

Note: PostgreSQL also has streaming replication slots (see Section 26.2.5), but they are used somewhat differently there.

A replication slot has an identifier that is unique across all databases in a PostgreSQL cluster. Slots persist independently of the connection using them and are crash-safe.

A logical slot will emit each change just once in normal operation. The current position of each slot is persisted only at checkpoint, so in the case of a crash the slot may return to an earlier LSN, which will then cause recent changes to be resent when the server restarts. Logical decoding clients are responsible for avoiding ill effects from handling the same message more than once. Clients may wish to record the last LSN they saw when decoding and skip over any repeated data or (when using the replication protocol) request that decoding start from that LSN rather than letting the server...
determine the start point. The Replication Progress Tracking feature is designed for this purpose, refer to replication origins.

Multiple independent slots may exist for a single database. Each slot has its own state, allowing different consumers to receive changes from different points in the database change stream. For most applications, a separate slot will be required for each consumer.

A logical replication slot knows nothing about the state of the receiver(s). It’s even possible to have multiple different receivers using the same slot at different times; they’ll just get the changes following on from when the last receiver stopped consuming them. Only one receiver may consume changes from a slot at any given time.

Note: Replication slots persist across crashes and know nothing about the state of their consumer(s). They will prevent removal of required resources even when there is no connection using them. This consumes storage because neither required WAL nor required rows from the system catalogs can be removed by VACUUM as long as they are required by a replication slot. So if a slot is no longer required it should be dropped.

47.2.3. Output Plugins

Output plugins transform the data from the write-ahead log’s internal representation into the format the consumer of a replication slot desires.

47.2.4. Exported Snapshots

When a new replication slot is created using the streaming replication interface, a snapshot is exported (see Section 9.26.5), which will show exactly the state of the database after which all changes will be included in the change stream. This can be used to create a new replica by using SET TRANSACTION SNAPSHOT to read the state of the database at the moment the slot was created. This transaction can then be used to dump the database’s state at that point in time, which afterwards can be updated using the slot’s contents without losing any changes.

47.3. Streaming Replication Protocol Interface

The commands

- CREATE_REPLICATION_SLOT slot_name LOGICAL output_plugin
- DROP_REPLICATION_SLOT slot_name
- START_REPLICATION SLOT slot_name LOGICAL ...

are used to create, drop, and stream changes from a replication slot, respectively. These commands are only available over a replication connection; they cannot be used via SQL. See Section 51.3 for details on these commands.

The command pg_recvlogical can be used to control logical decoding over a streaming replication connection. (It uses these commands internally.)
47.4. Logical Decoding SQL Interface

See Section 9.26.6 for detailed documentation on the SQL-level API for interacting with logical decoding.

Synchronous replication (see Section 26.2.8) is only supported on replication slots used over the streaming replication interface. The function interface and additional, non-core interfaces do not support synchronous replication.

47.5. System Catalogs Related to Logical Decoding

The `pg_replication_slots` view and the `pg_stat_replication` view provide information about the current state of replication slots and streaming replication connections respectively. These views apply to both physical and logical replication.

47.6. Logical Decoding Output Plugins

An example output plugin can be found in the `contrib/test_decoding` subdirectory of the PostgreSQL source tree.

47.6.1. Initialization Function

An output plugin is loaded by dynamically loading a shared library with the output plugin’s name as the library base name. The normal library search path is used to locate the library. To provide the required output plugin callbacks and to indicate that the library is actually an output plugin it needs to provide a function named `_PG_output_plugin_init`. This function is passed a struct that needs to be filled with the callback function pointers for individual actions.

```c
typedef struct OutputPluginCallbacks {
    LogicalDecodeStartupCB startup_cb;
    LogicalDecodeBeginCB begin_cb;
    LogicalDecodeChangeCB change_cb;
    LogicalDecodeCommitCB commit_cb;
    LogicalDecodeMessageCB message_cb;
    LogicalDecodeFilterByOriginCB filter_by_origin_cb;
    LogicalDecodeShutdownCB shutdown_cb;
} OutputPluginCallbacks;
```

```c
typedef void (*LogicalOutputPluginInit) (struct OutputPluginCallbacks *cb);
```

The `begin_cb`, `change_cb` and `commit_cb` callbacks are required, while `startup_cb`, `filter_by_origin_cb` and `shutdown_cb` are optional.

47.6.2. Capabilities

To decode, format and output changes, output plugins can use most of the backend’s normal infrastructure, including calling output functions. Read only access to relations is permitted as long as only
relations are accessed that either have been created by `initdb` in the `pg_catalog` schema, or have been marked as user provided catalog tables using

```sql
ALTER TABLE user_catalog_table SET (user_catalog_table = true);
CREATE TABLE another_catalog_table(data text) WITH (user_catalog_table = true);
```

Any actions leading to transaction ID assignment are prohibited. That, among others, includes writing to tables, performing DDL changes, and calling `txid_current()`.

### 47.6.3. Output Modes

Output plugin callbacks can pass data to the consumer in nearly arbitrary formats. For some use cases, like viewing the changes via SQL, returning data in a data type that can contain arbitrary data (e.g., `bytea`) is cumbersome. If the output plugin only outputs textual data in the server’s encoding, it can declare that by setting `OutputPluginOptions.output_type` to `OUTPUT_PLUGIN_TEXTUAL_OUTPUT` instead of `OUTPUT_PLUGIN_BINARY_OUTPUT` in the startup callback. In that case, all the data has to be in the server’s encoding so that a `text` datum can contain it. This is checked in assertion-enabled builds.

### 47.6.4. Output Plugin Callbacks

An output plugin gets notified about changes that are happening via various callbacks it needs to provide.

Concurrent transactions are decoded in commit order, and only changes belonging to a specific transaction are decoded between the `begin` and `commit` callbacks. Transactions that were rolled back explicitly or implicitly never get decoded. Successful savepoints are folded into the transaction containing them in the order they were executed within that transaction.

**Note:** Only transactions that have already safely been flushed to disk will be decoded. That can lead to a `COMMIT` not immediately being decoded in a directly following `pg_logical_slot_get_changes()` when `synchronous_commit` is set to `off`.

### 47.6.4.1. Startup Callback

The optional `startup_cb` callback is called whenever a replication slot is created or asked to stream changes, independent of the number of changes that are ready to be put out.

```c
typedef void (*LogicalDecodeStartupCB) (struct LogicalDecodingContext *ctx,
                                         OutputPluginOptions *options,
                                         bool is_init);
```

The `is_init` parameter will be true when the replication slot is being created and false otherwise. `options` points to a struct of options that output plugins can set:

```c
typedef struct OutputPluginOptions
{
    OutputPluginOutputType output_type;
} OutputPluginOptions;
```
output_type has to either be set to OUTPUT_PLUGIN_TEXTUAL_OUTPUT or OUTPUT_PLUGIN_BINARY_OUTPUT. See also Section 47.6.3.

The startup callback should validate the options present in ctx->output_plugin_options. If the output plugin needs to have a state, it can use ctx->output_plugin_private to store it.

### 47.6.4.2. Shutdown Callback

The optional shutdown_cb callback is called whenever a formerly active replication slot is not used anymore and can be used to deallocate resources private to the output plugin. The slot isn’t necessarily being dropped, streaming is just being stopped.

```c
typedef void (*LogicalDecodeShutdownCB) (struct LogicalDecodingContext *ctx);
```

### 47.6.4.3. Transaction Begin Callback

The required begin_cb callback is called whenever a start of a committed transaction has been decoded. Aborted transactions and their contents never get decoded.

```c
typedef void (*LogicalDecodeBeginCB) (struct LogicalDecodingContext *ctx,
    ReorderBufferTXN *txn);
```

The txn parameter contains meta information about the transaction, like the time stamp at which it has been committed and its XID.

### 47.6.4.4. Transaction End Callback

The required commit_cb callback is called whenever a transaction commit has been decoded. The change_cb callbacks for all modified rows will have been called before this, if there have been any modified rows.

```c
typedef void (*LogicalDecodeCommitCB) (struct LogicalDecodingContext *ctx,
    ReorderBufferTXN *txn,
    XLogRecPtr commit_lsn);
```

### 47.6.4.5. Change Callback

The required change_cb callback is called for every individual row modification inside a transaction, may it be an INSERT, UPDATE, or DELETE. Even if the original command modified several rows at once the callback will be called individually for each row.

```c
typedef void (*LogicalDecodeChangeCB) (struct LogicalDecodingContext *ctx,
    ReorderBufferTXN *txn,
    Relation relation,
    ReorderBufferChange *change);
```

The ctx and txn parameters have the same contents as for the begin_cb and commit_cb callbacks, but additionally the relation descriptor relation points to the relation the row belongs to and a struct change describing the row modification are passed in.
Chapter 47. Logical Decoding

Note: Only changes in user defined tables that are not unlogged (see UNLOGGED) and not temporary (see TEMPORARY or TEMP) can be extracted using logical decoding.

47.6.4.6. Origin Filter Callback

The optional filter_by_origin_cb callback is called to determine whether data that has been replayed from origin_id is of interest to the output plugin.

```c
typedef bool (*LogicalDecodeFilterByOriginCB) (struct LogicalDecodingContext *ctx, RepOriginId origin_id);
```

The ctx parameter has the same contents as for the other callbacks. No information but the origin is available. To signal that changes originating on the passed in node are irrelevant, return true, causing them to be filtered away; false otherwise. The other callbacks will not be called for transactions and changes that have been filtered away.

This is useful when implementing cascading or multidirectional replication solutions. Filtering by the origin allows to prevent replicating the same changes back and forth in such setups. While transactions and changes also carry information about the origin, filtering via this callback is noticeably more efficient.

47.6.4.7. Generic Message Callback

The optional message_cb callback is called whenever a logical decoding message has been decoded.

```c
typedef void (*LogicalDecodeMessageCB) (struct LogicalDecodingContext *ctx, ReorderBufferTXN *txn, XLogRecPtr message_lsn, bool transactional, const char *prefix, Size message_size, const char *message);
```

The txn parameter contains meta information about the transaction, like the time stamp at which it has been committed and its XID. Note however that it can be NULL when the message is non-transactional and the XID was not assigned yet in the transaction which logged the message. The lsn has WAL position of the message. The transactional says if the message was sent as transactional or not. The prefix is arbitrary null-terminated prefix which can be used for identifying interesting messages for the current plugin. And finally the message parameter holds the actual message of message_size size.

Extra care should be taken to ensure that the prefix the output plugin considers interesting is unique. Using name of the extension or the output plugin itself is often a good choice.

47.6.5. Functions for Producing Output

To actually produce output, output plugins can write data to the StringInfo output buffer in ctx->out when inside the begin_cb, commit_cb, or change_cb callbacks. Before writing to the output buffer, OutputPluginPrepareWrite(ctx, last_write) has to be called, and after finishing writing to the buffer, OutputPluginWrite(ctx, last_write) has to be called to
perform the write. The `last_write` indicates whether a particular write was the callback’s last write.

The following example shows how to output data to the consumer of an output plugin:

```c
OutputPluginPrepareWrite(ctx, true);
appendStringInfo(ctx->out, "BEGIN %u", txn->xid);
OutputPluginWrite(ctx, true);
```

### 47.7. Logical Decoding Output Writers

It is possible to add more output methods for logical decoding. For details, see the file `src/backend/replication/logical/logicalfuncs.c`. Essentially, three functions need to be provided: one to read WAL, one to prepare writing output, and one to write the output (see Section 47.6.5).

### 47.8. Synchronous Replication Support for Logical Decoding

Logical decoding can be used to build synchronous replication solutions with the same user interface as synchronous replication for streaming replication. To do this, the streaming replication interface (see Section 47.3) must be used to stream out data. Clients have to send `Standby status update` (F) (see Section 51.3) messages, just like streaming replication clients do.

**Note:** A synchronous replica receiving changes via logical decoding will work in the scope of a single database. Since, in contrast to that, `synchronous_standby_names` currently is server wide, this means this technique will not work properly if more than one database is actively used.
Chapter 48. Replication Progress Tracking

Replication origins are intended to make it easier to implement logical replication solutions on top of logical decoding. They provide a solution to two common problems:

- How to safely keep track of replication progress
- How to change replication behavior based on the origin of a row; for example, to prevent loops in bi-directional replication setups

Replication origins have just two properties, a name and an OID. The name, which is what should be used to refer to the origin across systems, is free-form text. It should be used in a way that makes conflicts between replication origins created by different replication solutions unlikely; e.g. by prefixing the replication solution’s name to it. The OID is used only to avoid having to store the long version in situations where space efficiency is important. It should never be shared across systems.

Replication origins can be created using the function `pg_replication_origin_create()`; dropped using `pg_replication_origin_drop()`; and seen in the `pg_replication_origin` system catalog.

One nontrivial part of building a replication solution is to keep track of replay progress in a safe manner. When the applying process, or the whole cluster, dies, it needs to be possible to find out up to where data has successfully been replicated. Naive solutions to this, such as updating a row in a table for every replayed transaction, have problems like run-time overhead and database bloat.

Using the replication origin infrastructure a session can be marked as replaying from a remote node (using the `pg_replication_origin_session_setup()` function). Additionally the LSN and commit time stamp of every source transaction can be configured on a per transaction basis using `pg_replication_origin_xact_setup()`. If that’s done replication progress will persist in a crash safe manner. Replay progress for all replication origins can be seen in the `pg_replication_origin_status` view. An individual origin’s progress, e.g. when resuming replication, can be acquired using `pg_replication_origin_progress()` for any origin or `pg_replication_origin_session_progress()` for the origin configured in the current session.

In replication topologies more complex than replication from exactly one system to one other system, another problem can be that it is hard to avoid replicating replayed rows again. That can lead both to cycles in the replication and inefficiencies. Replication origins provide an optional mechanism to recognize and prevent that. When configured using the functions referenced in the previous paragraph, every change and transaction passed to output plugin callbacks (see Section 47.6) generated by the session is tagged with the replication origin of the generating session. This allows treating them differently in the output plugin, e.g. ignoring all but locally-originating rows. Additionally the `filter_by_origin_cb` callback can be used to filter the logical decoding change stream based on the source. While less flexible, filtering via that callback is considerably more efficient than doing it in the output plugin.
VI. Reference

The entries in this Reference are meant to provide in reasonable length an authoritative, complete, and formal summary about their respective subjects. More information about the use of PostgreSQL, in narrative, tutorial, or example form, can be found in other parts of this book. See the cross-references listed on each reference page.

The reference entries are also available as traditional “man” pages.
I. SQL Commands

This part contains reference information for the SQL commands supported by PostgreSQL. By “SQL” the language in general is meant; information about the standards conformance and compatibility of each command can be found on the respective reference page.
ABORT

Name
ABORT — abort the current transaction

Synopsis
ABORT [ WORK | TRANSACTION ]

Description
ABORT rolls back the current transaction and causes all the updates made by the transaction to be discarded. This command is identical in behavior to the standard SQL command ROLLBACK, and is present only for historical reasons.

Parameters
WORK
TRANSACTION
Optional key words. They have no effect.

Notes
Use COMMIT to successfully terminate a transaction.
Issuing ABORT outside of a transaction block emits a warning and otherwise has no effect.

Examples
To abort all changes:

ABORT;

Compatibility
This command is a PostgreSQL extension present for historical reasons. ROLLBACK is the equivalent standard SQL command.
See Also

BEGIN, COMMIT, ROLLBACK
ALTER AGGREGATE

Name
ALTER AGGREGATE — change the definition of an aggregate function

Synopsis
ALTER AGGREGATE name ( aggregate_signature ) RENAME TO new_name
ALTER AGGREGATE name ( aggregate_signature )
    OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER AGGREGATE name ( aggregate_signature ) SET SCHEMA new_schema

where aggregate_signature is:

* | [ argmode ] [ argname ] argtype [ , ... ] |
[ [ argmode ] [ argname ] argtype [ , ... ] ] ORDER BY [ argmode ] [ argname ] argtype [ , ... ]

Description
ALTER AGGREGATE changes the definition of an aggregate function.

You must own the aggregate function to use ALTER AGGREGATE. To change the schema of an aggregate function, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the aggregate function’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the aggregate function. However, a superuser can alter ownership of any aggregate function anyway.)

Parameters

name
The name (optionally schema-qualified) of an existing aggregate function.

argmode
The mode of an argument: IN or VARIADIC. If omitted, the default is IN.

argname
The name of an argument. Note that ALTER AGGREGATE does not actually pay any attention to argument names, since only the argument data types are needed to determine the aggregate function’s identity.

argtype
An input data type on which the aggregate function operates. To reference a zero-argument aggregate function, write * in place of the list of argument specifications. To reference an ordered-set aggregate function, write ORDER BY between the direct and aggregated argument specifications.
new_name

The new name of the aggregate function.

new_owner

The new owner of the aggregate function.

new_schema

The new schema for the aggregate function.

Notes

The recommended syntax for referencing an ordered-set aggregate is to write `ORDER BY` between the direct and aggregated argument specifications, in the same style as in `CREATE AGGREGATE`. However, it will also work to omit `ORDER BY` and just run the direct and aggregated argument specifications into a single list. In this abbreviated form, if `VARIADIC "any"` was used in both the direct and aggregated argument lists, write `VARIADIC "any"` only once.

Examples

To rename the aggregate function `myavg` for `integer` to `my_average`:

```
ALTER AGGREGATE myavg(integer) RENAME TO my_average;
```

To change the owner of the aggregate function `myavg` for `integer` to `joe`:

```
ALTER AGGREGATE myavg(integer) OWNER TO joe;
```

To move the ordered-set aggregate `mypercentile` with direct argument of type `float8` and aggregated argument of type `integer` into schema `myschema`:

```
ALTER AGGREGATE mypercentile(float8 ORDER BY integer) SET SCHEMA myschema;
```
This will work too:

```
ALTER AGGREGATE mypercentile(float8, integer) SET SCHEMA myschema;
```

Compatibility

There is no `ALTER AGGREGATE` statement in the SQL standard.

See Also

`CREATE AGGREGATE`, `DROP AGGREGATE`
ALTER COLLATION

Name
ALTER COLLATION — change the definition of a collation

Synopsis
ALTER COLLATION  name RENAME TO  new_name
ALTER COLLATION  name OWNER TO  {  new_owner  |  CURRENT_USER  |  SESSION_USER  }
ALTER COLLATION  name SET SCHEMA  new_schema

Description
ALTER COLLATION changes the definition of a collation.
You must own the collation to use ALTER COLLATION. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the collation’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the collation. However, a superuser can alter ownership of any collation anyway.)

Parameters
name
The name (optionally schema-qualified) of an existing collation.
new_name
The new name of the collation.
new_owner
The new owner of the collation.
new_schema
The new schema for the collation.

Examples
To rename the collation de_DE to german:
ALTER COLLATION "de_DE" RENAME TO german;

To change the owner of the collation en_US to joe:
ALTER COLLATION "en_US" OWNER TO joe;
Compatibility

There is no ALTER COLLATION statement in the SQL standard.

See Also

CREATE COLLATION, DROP COLLATION
**ALTER CONVERSION**

**Name**
ALTER CONVERSION — change the definition of a conversion

**Synopsis**

ALTER CONVERSION name RENAME TO new_name  
ALTER CONVERSION name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }  
ALTER CONVERSION name SET SCHEMA new_schema

**Description**

ALTER CONVERSION changes the definition of a conversion.
You must own the conversion to use ALTER CONVERSION. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the conversion’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the conversion. However, a superuser can alter ownership of any conversion anyway.)

**Parameters**

name  
The name (optionally schema-qualified) of an existing conversion.  
new_name  
The new name of the conversion.  
new_owner  
The new owner of the conversion.  
new_schema  
The new schema for the conversion.

**Examples**

To rename the conversion *iso_8859_1_to_utf8* to *latin1_to_unicode*:

```
ALTER CONVERSION iso_8859_1_to_utf8 RENAME TO latin1_to_unicode;
```

To change the owner of the conversion *iso_8859_1_to_utf8* to *joe*:

```
ALTER CONVERSION iso_8859_1_to_utf8 OWNER TO joe;
```
Compatibility

There is no ALTER CONVERSION statement in the SQL standard.

See Also

CREATE CONVERSION, DROP CONVERSION
ALTER DATABASE

Name
ALTER DATABASE — change a database

Synopsis
ALTER DATABASE name [ [ WITH ] option [ ... ] ]

where option can be:

- ALLOW_CONNECTIONS allowconn
- CONNECTION LIMIT conlimit
- IS_TEMPLATE istemplate

ALTER DATABASE name RENAME TO new_name

ALTER DATABASE name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER DATABASE name SET TABLESPACE new_tablespace

ALTER DATABASE name SET configuration_parameter { TO | = } { value | DEFAULT }
ALTER DATABASE name SET configuration_parameter FROM CURRENT
ALTER DATABASE name RESET configuration_parameter
ALTER DATABASE name RESET ALL

Description
ALTER DATABASE changes the attributes of a database.

The first form changes certain per-database settings. (See below for details.) Only the database owner or a superuser can change these settings.

The second form changes the name of the database. Only the database owner or a superuser can rename a database; non-superuser owners must also have the CREATEDB privilege. The current database cannot be renamed. (Connect to a different database if you need to do that.)

The third form changes the owner of the database. To alter the owner, you must own the database and also be a direct or indirect member of the new owning role, and you must have the CREATEDB privilege. (Note that superusers have all these privileges automatically.)

The fourth form changes the default tablespace of the database. Only the database owner or a superuser can do this; you must also have create privilege for the new tablespace. This command physically moves any tables or indexes in the database’s old default tablespace to the new tablespace. The new default tablespace must be empty for this database, and no one can be connected to the database. Tables and indexes in non-default tablespaces are unaffected.

The remaining forms change the session default for a run-time configuration variable for a PostgreSQL database. Whenever a new session is subsequently started in that database, the specified value becomes the session default value. The database-specific default overrides whatever setting is present in postgresql.conf or has been received from the postgres command line. Only the database owner or a superuser can change the session defaults for a database. Certain variables cannot be set this way, or can only be set by a superuser.
Parameters

**name**

The name of the database whose attributes are to be altered.

**allowconn**

If false then no one can connect to this database.

**connlimit**

How many concurrent connections can be made to this database. -1 means no limit.

**istemplate**

If true, then this database can be cloned by any user with `CREATEDB` privileges; if false, then only superusers or the owner of the database can clone it.

**new_name**

The new name of the database.

**new_owner**

The new owner of the database.

**new_tablespace**

The new default tablespace of the database.

**configuration_parameter value**

Set this database’s session default for the specified configuration parameter to the given value. If `value` is `DEFAULT` or, equivalently, `RESET` is used, the database-specific setting is removed, so the system-wide default setting will be inherited in new sessions. Use `RESET ALL` to clear all database-specific settings. `SET FROM CURRENT` saves the session’s current value of the parameter as the database-specific value.

See SET and Chapter 19 for more information about allowed parameter names and values.

Notes

It is also possible to tie a session default to a specific role rather than to a database; see ALTER ROLE. Role-specific settings override database-specific ones if there is a conflict.

Examples

To disable index scans by default in the database `test`:

```
ALTER DATABASE test SET enable_indexscan TO off;
```

Compatibility

The `ALTER DATABASE` statement is a PostgreSQL extension.
See Also

CREATE DATABASE, DROP DATABASE, SET, CREATE TABLESPACE
ALTER DEFAULT PRIVILEGES

Name
ALTER DEFAULT PRIVILEGES — define default access privileges

Synopsis
ALTER DEFAULT PRIVILEGES
   [ FOR { ROLE | USER } target_role [, ...] ]
   [ IN SCHEMA schema_name [, ...] ]
abbreviated_grant_or_revoke

where abbreviated_grant_or_revoke is one of:

GRANT { { SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES | TRIGGER } [, ...] | ALL [ PRIVILEGES ] }
   ON TABLES
   TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

GRANT { { USAGE | SELECT | UPDATE } [, ...] | ALL [ PRIVILEGES ] }
   ON SEQUENCES
   TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

GRANT { EXECUTE | ALL [ PRIVILEGES ] }
   ON FUNCTIONS
   TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
   ON TYPES
   TO { [ GROUP ] role_name | PUBLIC } [, ...] [ WITH GRANT OPTION ]

REVOKE [ GRANT OPTION FOR ]
   { { SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES | TRIGGER } [, ...] | ALL [ PRIVILEGES ] }
   ON TABLES
   FROM { [ GROUP ] role_name | PUBLIC } [, ...]
   [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
   { { USAGE | SELECT | UPDATE } [, ...] | ALL [ PRIVILEGES ] }
   ON SEQUENCES
   FROM { [ GROUP ] role_name | PUBLIC } [, ...]
   [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
   { EXECUTE | ALL [ PRIVILEGES ] }
   ON FUNCTIONS
   FROM { [ GROUP ] role_name | PUBLIC } [, ...]
   [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
   { USAGE | ALL [ PRIVILEGES ] }

1306
ALTER DEFAULT PRIVILEGES

ON TYPES
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

Description

ALTER DEFAULT PRIVILEGES allows you to set the privileges that will be applied to objects created in the future. (It does not affect privileges assigned to already-existing objects.) Currently, only the privileges for tables (including views and foreign tables), sequences, functions, and types (including domains) can be altered.

You can change default privileges only for objects that will be created by yourself or by roles that you are a member of. The privileges can be set globally (i.e., for all objects created in the current database), or just for objects created in specified schemas.

As explained under GRANT, the default privileges for any object type normally grant all grantable permissions to the object owner, and may grant some privileges to PUBLIC as well. However, this behavior can be changed by altering the global default privileges with ALTER DEFAULT PRIVILEGES.

Default privileges that are specified per-schema are added to whatever the global default privileges are for the particular object type. This means you cannot revoke privileges per-schema if they are granted globally (either by default, or according to a previous ALTER DEFAULT PRIVILEGES command that did not specify a schema). Per-schema REVOKE is only useful to reverse the effects of a previous per-schema GRANT.

Parameters

target_role

The name of an existing role of which the current role is a member. If FOR ROLE is omitted, the current role is assumed.

schema_name

The name of an existing schema. If specified, the default privileges are altered for objects later created in that schema. If IN SCHEMA is omitted, the global default privileges are altered.

role_name

The name of an existing role to grant or revoke privileges for. This parameter, and all the other parameters in abbreviated_grant_or_revoke, act as described under GRANT or REVOKE, except that one is setting permissions for a whole class of objects rather than specific named objects.

Notes

Use psql’s \ddp command to obtain information about existing assignments of default privileges. The meaning of the privilege values is the same as explained for \dp under GRANT.

If you wish to drop a role for which the default privileges have been altered, it is necessary to reverse the changes in its default privileges or use DROP OWNED BY to get rid of the default privileges entry for the role.
Examples

Grant SELECT privilege to everyone for all tables (and views) you subsequently create in schema myschema, and allow role webuser to INSERT into them too:

```
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema GRANT SELECT ON TABLES TO PUBLIC;
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema GRANT INSERT ON TABLES TO webuser;
```

Undo the above, so that subsequently-created tables won’t have any more permissions than normal:

```
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema REVOKE SELECT ON TABLES FROM PUBLIC;
ALTER DEFAULT PRIVILEGES IN SCHEMA myschema REVOKE INSERT ON TABLES FROM webuser;
```

Remove the public EXECUTE permission that is normally granted on functions, for all functions subsequently created by role admin:

```
ALTER DEFAULT PRIVILEGES FOR ROLE admin REVOKE EXECUTE ON FUNCTIONS FROM PUBLIC;
```

Note however that you cannot accomplish that effect with a command limited to a single schema. This command has no effect, unless it is undoing a matching GRANT:

```
ALTER DEFAULT PRIVILEGES IN SCHEMA public REVOKE EXECUTE ON FUNCTIONS FROM PUBLIC;
```

That’s because per-schema default privileges can only add privileges to the global setting, not remove privileges granted by it.

Compatibility

There is no `ALTER DEFAULT PRIVILEGES` statement in the SQL standard.

See Also

`GRANT`, `REVOKE`
ALTER DOMAIN

Name
ALTER DOMAIN — change the definition of a domain

Synopsis
ALTER DOMAIN name
{ SET DEFAULT expression | DROP DEFAULT }
ALTER DOMAIN name
{ SET | DROP } NOT NULL
ALTER DOMAIN name
ADD domain_constraint [ NOT VALID ]
ALTER DOMAIN name
DROP CONSTRAINT [ IF EXISTS ] constraint_name [ RESTRICT | CASCADE ]
ALTER DOMAIN name
RENAME CONSTRAINT constraint_name TO new_constraint_name
ALTER DOMAIN name
VALIDATE CONSTRAINT constraint_name
ALTER DOMAIN name
OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER DOMAIN name
RENAME TO new_name
ALTER DOMAIN name
SET SCHEMA new_schema

Description
ALTER DOMAIN changes the definition of an existing domain. There are several sub-forms:

SET/DROP DEFAULT

These forms set or remove the default value for a domain. Note that defaults only apply to subsequent INSERT commands; they do not affect rows already in a table using the domain.

SET/DROP NOT NULL

These forms change whether a domain is marked to allow NULL values or to reject NULL values. You can only SET NOT NULL when the columns using the domain contain no null values.

ADD domain_constraint [ NOT VALID ]

This form adds a new constraint to a domain using the same syntax as CREATE DOMAIN. When a new constraint is added to a domain, all columns using that domain will be checked against the newly added constraint. These checks can be suppressed by adding the new constraint using the NOT VALID option; the constraint can later be made valid using ALTER DOMAIN ... VALIDATE CONSTRAINT. Newly inserted or updated rows are always checked against all constraints, even those marked NOT VALID. NOT VALID is only accepted for CHECK constraints.

DROP CONSTRAINT [ IF EXISTS ]

This form drops constraints on a domain. If IF EXISTS is specified and the constraint does not exist, no error is thrown. In this case a notice is issued instead.
RENAMe CONSTRAINT
   This form changes the name of a constraint on a domain.

VALIDATE CONSTRAINT
   This form validates a constraint previously added as NOT VALID, that is, it verifies that all values in table columns of the domain type satisfy the specified constraint.

OWNER
   This form changes the owner of the domain to the specified user.

RENAME
   This form changes the name of the domain.

SET SCHEMA
   This form changes the schema of the domain. Any constraints associated with the domain are moved into the new schema as well.

You must own the domain to use ALTER DOMAIN. To change the schema of a domain, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the domain’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the domain. However, a superuser can alter ownership of any domain anyway.)

Parameters

name
   The name (possibly schema-qualified) of an existing domain to alter.

domain_constraint
   New domain constraint for the domain.

constraint_name
   Name of an existing constraint to drop or rename.

NOT VALID
   Do not verify existing stored data for constraint validity.

CASCADE
   Automatically drop objects that depend on the constraint, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
   Refuse to drop the constraint if there are any dependent objects. This is the default behavior.

new_name
   The new name for the domain.

new_constraint_name
   The new name for the constraint.

new_owner
   The user name of the new owner of the domain.
ALTER DOMAIN

new_schema

The new schema for the domain.

Notes

Although ALTER DOMAIN ADD CONSTRAINT attempts to verify that existing stored data satisfies the new constraint, this check is not bulletproof, because the command cannot “see” table rows that are newly inserted or updated and not yet committed. If there is a hazard that concurrent operations might insert bad data, the way to proceed is to add the constraint using the NOT VALID option, commit that command, wait until all transactions started before that commit have finished, and then issue ALTER DOMAIN VALIDATE CONSTRAINT to search for data violating the constraint. This method is reliable because once the constraint is committed, all new transactions are guaranteed to enforce it against new values of the domain type.

Currently, ALTER DOMAIN ADD CONSTRAINT, ALTER DOMAIN VALIDATE CONSTRAINT, and ALTER DOMAIN SET NOT NULL will fail if the validated named domain or any derived domain is used within a composite-type column of any table in the database. They should eventually be improved to be able to verify the new constraint for such nested columns.

Examples

To add a NOT NULL constraint to a domain:

ALTER DOMAIN zipcode SET NOT NULL;

To remove a NOT NULL constraint from a domain:

ALTER DOMAIN zipcode DROP NOT NULL;

To add a check constraint to a domain:

ALTER DOMAIN zipcode ADD CONSTRAINT zipchk CHECK (char_length(VALUE) = 5);

To remove a check constraint from a domain:

ALTER DOMAIN zipcode DROP CONSTRAINT zipchk;

To rename a check constraint on a domain:

ALTER DOMAIN zipcode RENAME CONSTRAINT zipchk TO zip_check;

To move the domain into a different schema:

ALTER DOMAIN zipcode SET SCHEMA customers;
Compatibility

ALTER DOMAIN conforms to the SQL standard, except for the OWNER, RENAME, SET SCHEMA, and VALIDATE CONSTRAINT variants, which are PostgreSQL extensions. The NOT VALID clause of the ADD CONSTRAINT variant is also a PostgreSQL extension.

See Also

CREATE DOMAIN, DROP DOMAIN
ALTER EVENT TRIGGER

Name
ALTER EVENT TRIGGER — change the definition of an event trigger

Synopsis
ALTER EVENT TRIGGER name DISABLE
ALTER EVENT TRIGGER name ENABLE [ REPLICA | ALWAYS ]
ALTER EVENT TRIGGER name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER EVENT TRIGGER name RENAME TO new_name

Description
ALTER EVENT TRIGGER changes properties of an existing event trigger.
You must be superuser to alter an event trigger.

Parameters

name
The name of an existing trigger to alter.

new_owner
The user name of the new owner of the event trigger.

new_name
The new name of the event trigger.

DISABLE/ENABLE [ REPLICA | ALWAYS ] TRIGGER
These forms configure the firing of event triggers. A disabled trigger is still known to the system,
but is not executed when its triggering event occurs. See also session_replication_role.

Compatibility
There is no ALTER EVENT TRIGGER statement in the SQL standard.

See Also
CREATE EVENT TRIGGER, DROP EVENT TRIGGER
ALTER EXTENSION

Name
ALTER EXTENSION — change the definition of an extension

Synopsis
ALTER EXTENSION name UPDATE [ TO new_version ]
ALTER EXTENSION name SET SCHEMA new_schema
ALTER EXTENSION name ADD member_object
ALTER EXTENSION name DROP member_object

where member_object is:

ACCESS METHOD object_name |
AGGREGATE aggregate_name ( aggregate_signature ) |
CAST (source_type AS target_type) |
COLLATION object_name |
CONVERSION object_name |
DOMAIN object_name |
EVENT TRIGGER object_name |
FOREIGN DATA WRAPPER object_name |
FOREIGN TABLE object_name |
FUNCTION function_name ( [ [ argmode ] [ argname ] argtype [, ... ] ] ) |
MATERIALIZED VIEW object_name |
OPERATOR operator_name (left_type, right_type) |
OPERATOR CLASS object_name USING index_method |
OPERATOR FAMILY object_name USING index_method |
[ PROCEDURAL ] LANGUAGE object_name |
SCHEMA object_name |
SEQUENCE object_name |
SERVER object_name |
TABLE object_name |
TEXT SEARCH CONFIGURATION object_name |
TEXT SEARCH DICTIONARY object_name |
TEXT SEARCH PARSER object_name |
TEXT SEARCH TEMPLATE object_name |
TRANSFORM FOR type_name LANGUAGE lang_name |
TYPE object_name |
VIEW object_name |

and aggregate_signature is:

* |
[ argmode ] [ argname ] argtype [, ... ] |
[ [ argmode ] [ argname ] argtype [, ... ] ] ORDER BY [ argmode ] [ argname ] argtype [, ... ]
Description

ALTER EXTENSION changes the definition of an installed extension. There are several subforms:

**UPDATE**

This form updates the extension to a newer version. The extension must supply a suitable update script (or series of scripts) that can modify the currently-installed version into the requested version.

**SET SCHEMA**

This form moves the extension’s objects into another schema. The extension has to be *relocatable* for this command to succeed.

**ADD member_object**

This form adds an existing object to the extension. This is mainly useful in extension update scripts. The object will subsequently be treated as a member of the extension; notably, it can only be dropped by dropping the extension.

**DROP member_object**

This form removes a member object from the extension. This is mainly useful in extension update scripts. The object is not dropped, only disassociated from the extension.

See Section 36.15 for more information about these operations.

You must own the extension to use `ALTER EXTENSION`. The `ADD/DROP` forms require ownership of the added/dropped object as well.

Parameters

- **name**
  
  The name of an installed extension.

- **new_version**

  The desired new version of the extension. This can be written as either an identifier or a string literal. If not specified, `ALTER EXTENSION UPDATE` attempts to update to whatever is shown as the default version in the extension’s control file.

- **new_schema**

  The new schema for the extension.

- **object_name**
  **aggregate_name**
  **function_name**
  **operator_name**

  The name of an object to be added to or removed from the extension. Names of tables, aggregates, domains, foreign tables, functions, operators, operator classes, operator families, sequences, text search objects, types, and views can be schema-qualified.

- **source_type**

  The name of the source data type of the cast.
**ALTER EXTENSION**

`target_type`

The name of the target data type of the cast.

`argmode`

The mode of a function or aggregate argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Note that `ALTER EXTENSION` does not actually pay any attention to OUT arguments, since only the input arguments are needed to determine the function’s identity. So it is sufficient to list the IN, INOUT, and VARIADIC arguments.

`argname`

The name of a function or aggregate argument. Note that `ALTER EXTENSION` does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

`argtype`

The data type of a function or aggregate argument.

`left_type`

`right_type`

The data type(s) of the operator’s arguments (optionally schema-qualified). Write NONE for the missing argument of a prefix or postfix operator.

**PROCEDURAL**

This is a noise word.

`type_name`

The name of the data type of the transform.

`lang_name`

The name of the language of the transform.

**Examples**

To update the `hstore` extension to version 2.0:

```
ALTER EXTENSION hstore UPDATE TO '2.0';
```

To change the schema of the `hstore` extension to `utils`:

```
ALTER EXTENSION hstore SET SCHEMA utils;
```

To add an existing function to the `hstore` extension:

```
ALTER EXTENSION hstore ADD FUNCTION populate_record(anyelement, hstore);
```
Compatibility

ALTER EXTENSION is a PostgreSQL extension.

See Also

CREATE EXTENSION, DROP EXTENSION
ALTER FOREIGN DATA WRAPPER

Name

ALTER FOREIGN DATA WRAPPER — change the definition of a foreign-data wrapper

Synopsis

ALTER FOREIGN DATA WRAPPER name
   [ HANDLER handler_function | NO HANDLER ]
   [ VALIDATOR validator_function | NO VALIDATOR ]
   [ OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] ) ]
ALTER FOREIGN DATA WRAPPER name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }  
ALTER FOREIGN DATA WRAPPER name RENAME TO new_name

Description

ALTER FOREIGN DATA WRAPPER changes the definition of a foreign-data wrapper. The first form of the command changes the support functions or the generic options of the foreign-data wrapper (at least one clause is required). The second form changes the owner of the foreign-data wrapper.

Only superusers can alter foreign-data wrappers. Additionally, only superusers can own foreign-data wrappers.

Parameters

name
   The name of an existing foreign-data wrapper.

HANDLER handler_function
   Specifies a new handler function for the foreign-data wrapper.

NO HANDLER
   This is used to specify that the foreign-data wrapper should no longer have a handler function.

   Note that foreign tables that use a foreign-data wrapper with no handler cannot be accessed.

VALIDATOR validator_function
   Specifies a new validator function for the foreign-data wrapper.

   Note that it is possible that pre-existing options of the foreign-data wrapper, or of dependent servers, user mappings, or foreign tables, are invalid according to the new validator. PostgreSQL does not check for this. It is up to the user to make sure that these options are correct before using the modified foreign-data wrapper. However, any options specified in this ALTER FOREIGN DATA WRAPPER command will be checked using the new validator.

NO VALIDATOR
   This is used to specify that the foreign-data wrapper should no longer have a validator function.
OPTIONS ( [ ADD | SET | DROP ] option [‘value’] [, ... ] )

Change options for the foreign-data wrapper. ADD, SET, and DROP specify the action to be performed. ADD is assumed if no operation is explicitly specified. Option names must be unique; names and values are also validated using the foreign data wrapper’s validator function, if any.

new_owner

The user name of the new owner of the foreign-data wrapper.

new_name

The new name for the foreign-data wrapper.

Examples

Change a foreign-data wrapper dbi, add option foo, drop bar:

```
ALTER FOREIGN DATA WRAPPER dbi OPTIONS (ADD foo ’1’, DROP ’bar’);
```

Change the foreign-data wrapper dbi validator to bob.myvalidator:

```
ALTER FOREIGN DATA WRAPPER dbi VALIDATOR bob.myvalidator;
```

Compatibility

ALTER FOREIGN DATA WRAPPER conforms to ISO/IEC 9075-9 (SQL/MED), except that the HANDLER, VALIDATOR, OWNER TO, and RENAME clauses are extensions.

See Also

CREATE FOREIGN DATA WRAPPER, DROP FOREIGN DATA WRAPPER
ALTER FOREIGN TABLE

Name
ALTER FOREIGN TABLE — change the definition of a foreign table

Synopsis

ALTER FOREIGN TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
   action [, ... ]
ALTER FOREIGN TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
   RENAME [ COLUMN ] column_name TO new_column_name
ALTER FOREIGN TABLE [ IF EXISTS ] name
   RENAME TO new_name
ALTER FOREIGN TABLE [ IF EXISTS ] name
   SET SCHEMA new_schema

where action is one of:

ADD [ COLUMN ] column_name data_type [ COLLATE collation ] [ column_constraint [ ... ] ]
DROP [ COLUMN ] [ IF EXISTS ] column_name [ RESTRICT | CASCADE ]
ALTER [ COLUMN ] column_name [ SET DATA ] TYPE data_type [ COLLATE collation ]
ALTER [ COLUMN ] column_name SET DEFAULT expression
ALTER [ COLUMN ] column_name DROP DEFAULT
ALTER [ COLUMN ] column_name { SET | DROP } NOT NULL
ALTER [ COLUMN ] column_name SET STATISTICS integer
ALTER [ COLUMN ] column_name SET ( attribute_option = value [, ... ] )
ALTER [ COLUMN ] column_name RESET ( attribute_option [, ... ] )
ALTER [ COLUMN ] column_name SET STORAGE { PLAIN | EXTERNAL | EXTENDED | MAIN }
ALTER [ COLUMN ] column_name OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ])
ADD table_constraint [ NOT VALID ]
VALIDATE CONSTRAINT constraint_name
DROP CONSTRAINT [ IF EXISTS ] constraint_name [ RESTRICT | CASCADE ]
DISABLE TRIGGER [ trigger_name | ALL | USER ]
ENABLE TRIGGER [ trigger_name | ALL | USER ]
ENABLE REPLICA TRIGGER trigger_name
ENABLE ALWAYS TRIGGER trigger_name
SET WITH OIDS
SET WITHOUT OIDS
INHERIT parent_table
NO INHERIT parent_table
OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ])

Description

ALTER FOREIGN TABLE changes the definition of an existing foreign table. There are several sub-forms:

ADD COLUMN

This form adds a new column to the foreign table, using the same syntax as CREATE FOREIGN
TABLE. Unlike the case when adding a column to a regular table, nothing happens to the under-
ALTER FOREIGN TABLE

lying storage; this action simply declares that some new column is now accessible through the foreign table.

DROP COLUMN [ IF EXISTS ]

This form drops a column from a foreign table. You will need to say CASCADE if anything outside the table depends on the column; for example, views. If IF EXISTS is specified and the column does not exist, no error is thrown. In this case a notice is issued instead.

SET DATA TYPE

This form changes the type of a column of a foreign table. Again, this has no effect on any underlying storage: this action simply changes the type that PostgreSQL believes the column to have.

SET/DROP DEFAULT

These forms set or remove the default value for a column. Default values only apply in subsequent INSERT or UPDATE commands; they do not cause rows already in the table to change.

SET/DROP NOT NULL

Mark a column as allowing, or not allowing, null values.

SET STATISTICS

This form sets the per-column statistics-gathering target for subsequent ANALYZE operations. See the similar form of ALTER TABLE for more details.

SET ( attribute_option = value [, ... ] )
RESET ( attribute_option [, ... ] )

This form sets or resets per-attribute options. See the similar form of ALTER TABLE for more details.

SET STORAGE

This form sets the storage mode for a column. See the similar form of ALTER TABLE for more details. Note that the storage mode has no effect unless the table’s foreign-data wrapper chooses to pay attention to it.

ADD table_constraint [ NOT VALID ]

This form adds a new constraint to a foreign table, using the same syntax as CREATE FOREIGN TABLE. Currently only CHECK constraints are supported.

Unlike the case when adding a constraint to a regular table, nothing is done to verify the constraint is correct; rather, this action simply declares that some new condition should be assumed to hold for all rows in the foreign table. (See the discussion in CREATE FOREIGN TABLE.) If the constraint is marked NOT VALID, then it isn’t assumed to hold, but is only recorded for possible future use.

VALIDATE CONSTRAINT

This form marks as valid a constraint that was previously marked as NOT VALID. No action is taken to verify the constraint, but future queries will assume that it holds.

DROP CONSTRAINT [ IF EXISTS ]

This form drops the specified constraint on a foreign table. If IF EXISTS is specified and the constraint does not exist, no error is thrown. In this case a notice is issued instead.

DISABLE/ENABLE [ REPLICA | ALWAYS ] TRIGGER

These forms configure the firing of trigger(s) belonging to the foreign table. See the similar form of ALTER TABLE for more details.
ALTER FOREIGN TABLE

SET WITH OIDS

This form adds an oid system column to the table (see Section 5.4). It does nothing if the table already has OIDs. Unless the table’s foreign-data wrapper supports OIDs, this column will simply read as zeroes.

Note that this is not equivalent to ADD COLUMN oid oid; that would add a normal column that happened to be named oid, not a system column.

SET WITHOUT OIDS

This form removes the oid system column from the table. This is exactly equivalent to DROP COLUMN oid RESTRICT, except that it will not complain if there is already no oid column.

INHERIT parent_table

This form adds the target foreign table as a new child of the specified parent table. See the similar form of ALTER TABLE for more details.

NO INHERIT parent_table

This form removes the target foreign table from the list of children of the specified parent table.

OWNER

This form changes the owner of the foreign table to the specified user.

OPTIONS ( [ ADD | SET | DROP ] option [‘value’] [, ... ] )

Change options for the foreign table or one of its columns. ADD, SET, and DROP specify the action to be performed. ADD is assumed if no operation is explicitly specified. Duplicate option names are not allowed (although it’s OK for a table option and a column option to have the same name).

Option names and values are also validated using the foreign data wrapper library.

RENAME

The RENAME forms change the name of a foreign table or the name of an individual column in a foreign table.

SET SCHEMA

This form moves the foreign table into another schema.

All the actions except RENAME and SET SCHEMA can be combined into a list of multiple alterations to apply in parallel. For example, it is possible to add several columns and/or alter the type of several columns in a single command.

If the command is written as ALTER FOREIGN TABLE IF EXISTS ... and the foreign table does not exist, no error is thrown. A notice is issued in this case.

You must own the table to use ALTER FOREIGN TABLE. To change the schema of a foreign table, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the table’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the table. However, a superuser can alter ownership of any table anyway.) To add a column or alter a column type, you must also have USAGE privilege on the data type.
ALTER FOREIGN TABLE

Parameters

name
The name (possibly schema-qualified) of an existing foreign table to alter. If ONLY is specified before the table name, only that table is altered. If ONLY is not specified, the table and all its descendant tables (if any) are altered. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

column_name
Name of a new or existing column.

new_column_name
New name for an existing column.

new_name
New name for the table.

data_type
Data type of the new column, or new data type for an existing column.

table_constraint
New table constraint for the foreign table.

constraint_name
Name of an existing constraint to drop.

CASCADE
Automatically drop objects that depend on the dropped column or constraint (for example, views referencing the column), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the column or constraint if there are any dependent objects. This is the default behavior.

trigger_name
Name of a single trigger to disable or enable.

ALL
Disable or enable all triggers belonging to the foreign table. (This requires superuser privilege if any of the triggers are internally generated triggers. The core system does not add such triggers to foreign tables, but add-on code could do so.)

USER
Disable or enable all triggers belonging to the foreign table except for internally generated triggers.

parent_table
A parent table to associate or de-associate with this foreign table.

new_owner
The user name of the new owner of the table.
new_schema

The name of the schema to which the table will be moved.

Notes

The key word COLUMN is noise and can be omitted.

Consistency with the foreign server is not checked when a column is added or removed with ADD COLUMN or DROP COLUMN, a NOT NULL or CHECK constraint is added, or a column type is changed with SET DATA TYPE. It is the user’s responsibility to ensure that the table definition matches the remote side.

Refer to CREATE FOREIGN TABLE for a further description of valid parameters.

Examples

To mark a column as not-null:

ALTER FOREIGN TABLE distributors ALTER COLUMN street SET NOT NULL;

To change options of a foreign table:

ALTER FOREIGN TABLE myschema.distributors OPTIONS (ADD opt1 ‘value’, SET opt2 ‘value2’, DROP opt3 ‘value3’);

Compatibility

The forms ADD, DROP, and SET DATA TYPE conform with the SQL standard. The other forms are PostgreSQL extensions of the SQL standard. Also, the ability to specify more than one manipulation in a single ALTER FOREIGN TABLE command is an extension.

ALTER FOREIGN TABLE DROP COLUMN can be used to drop the only column of a foreign table, leaving a zero-column table. This is an extension of SQL, which disallows zero-column foreign tables.

See Also

CREATE FOREIGN TABLE, DROP FOREIGN TABLE
**ALTER FUNCTION**

**Name**

ALTER FUNCTION — change the definition of a function

**Synopsis**

ALTER FUNCTION name ( [ [ argmode ] [ argname ] argtype [, ...] ] )

   action [ ... ] [ RESTRICT ]

ALTER FUNCTION name ( [ [ argmode ] [ argname ] argtype [, ...] ] )

   RENAME TO new_name

ALTER FUNCTION name ( [ [ argmode ] [ argname ] argtype [, ...] ] )

   OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER FUNCTION name ( [ [ argmode ] [ argname ] argtype [, ...] ] )

   SET SCHEMA new_schema

ALTER FUNCTION name ( [ [ argmode ] [ argname ] argtype [, ...] ] )

   DEPENDS ON EXTENSION extension_name

where action is one of:

   CALLED ON NULL INPUT | RETURNS NULL ON NULL INPUT | STRICT
   IMMUTABLE | STABLE | VOLATILE | [ NOT ] LEAKPROOF
   [ EXTERNAL ] SECURITY INVOKER | [ EXTERNAL ] SECURITY DEFINER
   PARALLEL { UNSAFE | RESTRICTED | SAFE }
   COST execution_cost
   ROWS result_rows
   SET configuration_parameter { TO | = } { value | DEFAULT }
   SET configuration_parameter FROM CURRENT
   RESET configuration_parameter
   RESET ALL

**Description**

ALTER FUNCTION changes the definition of a function.

You must own the function to use ALTER FUNCTION. To change a function’s schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the function’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the function. However, a superuser can alter ownership of any function anyway.)

**Parameters**

name

The name (optionally schema-qualified) of an existing function.
ALTER FUNCTION

argmode
The mode of an argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Note that ALTER FUNCTION does not actually pay any attention to OUT arguments, since only the input arguments are needed to determine the function’s identity. So it is sufficient to list the IN, INOUT, and VARIADIC arguments.

argname
The name of an argument. Note that ALTER FUNCTION does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

argtype
The data type(s) of the function’s arguments (optionally schema-qualified), if any.

new_name
The new name of the function.

new_owner
The new owner of the function. Note that if the function is marked SECURITY DEFINER, it will subsequently execute as the new owner.

new_schema
The new schema for the function.

extension_name
The name of the extension that the function is to depend on.

CALLED ON NULL INPUT
RETURNS NULL ON NULL INPUT
STRICT

CALLED ON NULL INPUT changes the function so that it will be invoked when some or all of its arguments are null. RETURNS NULL ON NULL INPUT or STRICT changes the function so that it is not invoked if any of its arguments are null; instead, a null result is assumed automatically. See CREATE FUNCTION for more information.

IMMUTABLE
STABLE
VOLATILE

Change the volatility of the function to the specified setting. See CREATE FUNCTION for details.

[ EXTERNAL ] SECURITY INVOKER
[ EXTERNAL ] SECURITY DEFINER

Change whether the function is a security definer or not. The key word EXTERNAL is ignored for SQL conformance. See CREATE FUNCTION for more information about this capability.

PARALLEL
Change whether the function is deemed safe for parallelism. See CREATE FUNCTION for details.

LEAKPROOF
Change whether the function is considered leakproof or not. See CREATE FUNCTION for more information about this capability.
COST execution_cost

Change the estimated execution cost of the function. See CREATE FUNCTION for more information.

ROWS result_rows

Change the estimated number of rows returned by a set-returning function. See CREATE FUNCTION for more information.

configuration_parameter value

Add or change the assignment to be made to a configuration parameter when the function is called. If \texttt{value} is \texttt{DEFAULT} or, equivalently, \texttt{RESET} is used, the function-local setting is removed, so that the function executes with the value present in its environment. Use \texttt{RESET ALL} to clear all function-local settings. \texttt{SET FROM CURRENT} saves the value of the parameter that is current when \texttt{ALTER FUNCTION} is executed as the value to be applied when the function is entered.

See \texttt{SET} and Chapter 19 for more information about allowed parameter names and values.

RESTRICT

Ignored for conformance with the SQL standard.

Examples

To rename the function \texttt{sqrt} for type \texttt{integer} to \texttt{square_root}:

\[
\text{ALTER FUNCTION sqrt(integer) RENAME TO square_root;}
\]

To change the owner of the function \texttt{sqrt} for type \texttt{integer} to \texttt{joe}:

\[
\text{ALTER FUNCTION sqrt(integer) OWNER TO joe;}
\]

To change the schema of the function \texttt{sqrt} for type \texttt{integer} to \texttt{maths}:

\[
\text{ALTER FUNCTION sqrt(integer) SET SCHEMA maths;}
\]

To mark the function \texttt{sqrt} for type \texttt{integer} as being dependent on the extension \texttt{mathlib}:

\[
\text{ALTER FUNCTION sqrt(integer) DEPENDS ON EXTENSION mathlib;}
\]

To adjust the search path that is automatically set for a function:

\[
\text{ALTER FUNCTION check_password(text) SET search_path = admin, pg_temp;}
\]

To disable automatic setting of \texttt{search_path} for a function:

\[
\text{ALTER FUNCTION check_password(text) RESET search_path;}
\]
ALTER FUNCTION

The function will now execute with whatever search path is used by its caller.

Compatibility

This statement is partially compatible with the ALTER FUNCTION statement in the SQL standard. The standard allows more properties of a function to be modified, but does not provide the ability to rename a function, make a function a security definer, attach configuration parameter values to a function, or change the owner, schema, or volatility of a function. The standard also requires the RESTRICT key word, which is optional in PostgreSQL.

See Also

CREATE FUNCTION, DROP FUNCTION
**ALTER GROUP**

**Name**
ALTER GROUP — change role name or membership

**Synopsis**

ALTER GROUP role_specification ADD USER user_name [, ... ]
ALTER GROUP role_specification DROP USER user_name [, ... ]

where role_specification can be:

- role_name
- CURRENT_USER
- SESSION_USER

ALTER GROUP group_name RENAME TO new_name

**Description**

ALTER GROUP changes the attributes of a user group. This is an obsolete command, though still accepted for backwards compatibility, because groups (and users too) have been superseded by the more general concept of roles.

The first two variants add users to a group or remove them from a group. (Any role can play the part of either a “user” or a “group” for this purpose.) These variants are effectively equivalent to granting or revoking membership in the role named as the “group”; so the preferred way to do this is to use GRANT or REVOKE.

The third variant changes the name of the group. This is exactly equivalent to renaming the role with ALTER ROLE.

**Parameters**

- **group_name**
  The name of the group (role) to modify.

- **user_name**
  Users (roles) that are to be added to or removed from the group. The users must already exist; ALTER GROUP does not create or drop users.

- **new_name**
  The new name of the group.
Examples

Add users to a group:

ALTER GROUP staff ADD USER karl, john;

Remove a user from a group:

ALTER GROUP workers DROP USER beth;

Compatibility

There is no ALTER GROUP statement in the SQL standard.

See Also

GRANT, REVOKE, ALTER ROLE
ALTER INDEX

Name
ALTER INDEX — change the definition of an index

Synopsis
ALTER INDEX [ IF EXISTS ] name RENAME TO new_name
ALTER INDEX [ IF EXISTS ] name SET TABLESPACE tablespace_name
ALTER INDEX name DEPENDS ON EXTENSION extension_name
ALTER INDEX [ IF EXISTS ] name SET { storage_parameter = value [, ... ] }
ALTER INDEX [ IF EXISTS ] name RESET { storage_parameter [, ... ] }
ALTER INDEX ALL IN TABLESPACE name [ OWNED BY role_name [, ... ] ]

SET TABLESPACE new_tablespace [ NOWAIT ]

Description
ALTER INDEX changes the definition of an existing index. There are several subforms:

RENAME
The RENAME form changes the name of the index. There is no effect on the stored data.

SET TABLESPACE
This form changes the index’s tablespace to the specified tablespace and moves the data file(s) associated with the index to the new tablespace. To change the tablespace of an index, you must own the index and have CREATE privilege on the new tablespace. All indexes in the current database in a tablespace can be moved by using the ALL IN TABLESPACE form, which will lock all indexes to be moved and then move each one. This form also supports OWNED BY, which will only move indexes owned by the roles specified. If the NOWAIT option is specified then the command will fail if it is unable to acquire all of the locks required immediately. Note that system catalogs will not be moved by this command, use ALTER DATABASE or explicit ALTER INDEX invocations instead if desired. See also CREATE TABLESPACE.

DEPENDS ON EXTENSION
This form marks the index as dependent on the extension, such that if the extension is dropped, the index will automatically be dropped as well.

SET { storage_parameter = value [, ... ] }
This form changes one or more index-method-specific storage parameters for the index. See CREATE INDEX for details on the available parameters. Note that the index contents will not be modified immediately by this command; depending on the parameter you might need to rebuild the index with REINDEX to get the desired effects.

RESET { storage_parameter [, ... ] }
This form resets one or more index-method-specific storage parameters to their defaults. As with SET, a REINDEX might be needed to update the index entirely.
Parameters

IF EXISTS
Do not throw an error if the index does not exist. A notice is issued in this case.

name
The name (possibly schema-qualified) of an existing index to alter.

new_name
The new name for the index.

tablespace_name
The tablespace to which the index will be moved.

extension_name
The name of the extension that the index is to depend on.

storage_parameter
The name of an index-method-specific storage parameter.

value
The new value for an index-method-specific storage parameter. This might be a number or a word depending on the parameter.

Notes
These operations are also possible using ALTER TABLE. ALTER INDEX is in fact just an alias for the forms of ALTER TABLE that apply to indexes.

There was formerly an ALTER INDEX OWNER variant, but this is now ignored (with a warning). An index cannot have an owner different from its table’s owner. Changing the table’s owner automatically changes the index as well.

Changing any part of a system catalog index is not permitted.

Examples
To rename an existing index:

ALTER INDEX distributors RENAME TO suppliers;

To move an index to a different tablespace:

ALTER INDEX distributors SET TABLESPACE fasttablespace;

To change an index’s fill factor (assuming that the index method supports it):

ALTER INDEX distributors SET (fillfactor = 75);
REINDEX INDEX distributors;
Compatibility

 ALTER INDEX is a PostgreSQL extension.

See Also

 CREATE INDEX, REINDEX
ALTER LANGUAGE

Name
ALTER LANGUAGE — change the definition of a procedural language

Synopsis
ALTER [ PROCEDURAL ] LANGUAGE name RENAME TO new_name
ALTER [ PROCEDURAL ] LANGUAGE name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description
ALTER LANGUAGE changes the definition of a procedural language. The only functionality is to rename the language or assign a new owner. You must be superuser or owner of the language to use ALTER LANGUAGE.

Parameters
name
Name of a language
new_name
The new name of the language
new_owner
The new owner of the language

Compatibility
There is no ALTER LANGUAGE statement in the SQL standard.

See Also
CREATE LANGUAGE, DROP LANGUAGE
ALTER LARGE OBJECT

Name
ALTER LARGE OBJECT — change the definition of a large object

Synopsis
ALTER LARGE OBJECT large_object_oid OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description
ALTER LARGE OBJECT changes the definition of a large object.
You must own the large object to use ALTER LARGE OBJECT. To alter the owner, you must also be
a direct or indirect member of the new owning role. (However, a superuser can alter any large object
anyway.) Currently, the only functionality is to assign a new owner, so both restrictions always apply.

Parameters
large_object_oid
OID of the large object to be altered
new_owner
The new owner of the large object

Compatibility
There is no ALTER LARGE OBJECT statement in the SQL standard.

See Also
Chapter 33
ALTER MATERIALIZED VIEW

Name

ALTER MATERIALIZED VIEW — change the definition of a materialized view

Synopsis

ALTER MATERIALIZED VIEW [ IF EXISTS ] name
    action [, ... ]
ALTER MATERIALIZED VIEW name
    DEPENDS ON EXTENSION extension_name
ALTER MATERIALIZED VIEW [ IF EXISTS ] name
    RENAME [ COLUMN ] column_name TO new_column_name
ALTER MATERIALIZED VIEW [ IF EXISTS ] name
    RENAME TO new_name
ALTER MATERIALIZED VIEW [ IF EXISTS ] name
    SET SCHEMA new_schema
ALTER MATERIALIZED VIEW ALL IN TABLESPACE name [ OWNED BY role_name [, ... ] ]
    SET TABLESPACE new_tablespace [ NOWAIT ]

where action is one of:

    ALTER [ COLUMN ] column_name SET STATISTICS integer
    ALTER [ COLUMN ] column_name SET ( attribute_option = value [, ... ] )
    ALTER [ COLUMN ] column_name RESET ( attribute_option [, ... ] )
    ALTER [ COLUMN ] column_name SET STORAGE { PLAIN | EXTERNAL | EXTENDED | MAIN }
    CLUSTER ON index_name
    SET WITHOUT CLUSTER
    SET ( storage_parameter = value [, ... ] )
    RESET ( storage_parameter [, ... ] )
    OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description

ALTER MATERIALIZED VIEW changes various auxiliary properties of an existing materialized view.

You must own the materialized view to use ALTER MATERIALIZED VIEW. To change a materialized view’s schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the materialized view’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the materialized view. However, a superuser can alter ownership of any view anyway.)

The DEPENDS ON EXTENSION form marks the materialized view as dependent on an extension, such that the materialized view will automatically be dropped if the extension is dropped.

The statement subforms and actions available for ALTER MATERIALIZED VIEW are a subset of those available for ALTER TABLE, and have the same meaning when used for materialized views. See the descriptions for ALTER TABLE for details.
ALTER MATERIALIZED VIEW

Parameters

name
The name (optionally schema-qualified) of an existing materialized view.

column_name
Name of a new or existing column.

extension_name
The name of the extension that the materialized view is to depend on.

new_column_name
New name for an existing column.

new_owner
The user name of the new owner of the materialized view.

new_name
The new name for the materialized view.

new_schema
The new schema for the materialized view.

Examples
To rename the materialized view foo to bar:

ALTER MATERIALIZED VIEW foo RENAME TO bar;

Compatibility

ALTER MATERIALIZED VIEW is a PostgreSQL extension.

See Also

CREATE MATERIALIZED VIEW, DROP MATERIALIZED VIEW, REFRESH MATERIALIZED VIEW
ALTER OPERATOR

Name
ALTER OPERATOR — change the definition of an operator

Synopsis
ALTER OPERATOR name { { left_type | NONE } , { right_type | NONE } } 
   OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER OPERATOR name { { left_type | NONE } , { right_type | NONE } } 
   SET SCHEMA new_schema

ALTER OPERATOR name { { left_type | NONE } , { right_type | NONE } } 
   SET { { RESTRICT = { res_proc | NONE } 
         | JOIN = { join_proc | NONE } 
       } [, ... ] }

Description
ALTER OPERATOR changes the definition of an operator.
You must own the operator to use ALTER OPERATOR. To alter the owner, you must also be a direct
or indirect member of the new owning role, and that role must have CREATE privilege on the opera-
tor’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do
by dropping and recreating the operator. However, a superuser can alter ownership of any operator
anyway.)

Parameters

name
   The name (optionally schema-qualified) of an existing operator.

left_type
   The data type of the operator’s left operand; write NONE if the operator has no left operand.

right_type
   The data type of the operator’s right operand; write NONE if the operator has no right operand.

new_owner
   The new owner of the operator.

new_schema
   The new schema for the operator.

res_proc
   The restriction selectivity estimator function for this operator; write NONE to remove existing
   selectivity estimator.
**ALTER OPERATOR**

join_proc

The join selectivity estimator function for this operator; write NONE to remove existing selectivity estimator.

**Examples**

Change the owner of a custom operator `a @@ b` for type `text`:

```sql
ALTER OPERATOR @@ (text, text) OWNER TO joe;
```

Change the restriction and join selectivity estimator functions of a custom operator `a && b` for type `int[]`:

```sql
ALTER OPERATOR && (_int4, _int4) SET (RESTRICT = _int_contsel, JOIN = _int_contjoinsel);
```

**Compatibility**

There is no `ALTER OPERATOR` statement in the SQL standard.

**See Also**

`CREATE OPERATOR`, `DROP OPERATOR`
ALTER OPERATOR CLASS

Name

ALTER OPERATOR CLASS — change the definition of an operator class

Synopsis

ALTER OPERATOR CLASS name USING index_method
    RENAME TO new_name

ALTER OPERATOR CLASS name USING index_method
    OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER OPERATOR CLASS name USING index_method
    SET SCHEMA new_schema

Description

ALTER OPERATOR CLASS changes the definition of an operator class.

You must own the operator class to use ALTER OPERATOR CLASS. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the operator class’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the operator class. However, a superuser can alter ownership of any operator class anyway.)

Parameters

name

The name (optionally schema-qualified) of an existing operator class.

index_method

The name of the index method this operator class is for.

new_name

The new name of the operator class.

new_owner

The new owner of the operator class.

new_schema

The new schema for the operator class.

Compatibility

There is no ALTER OPERATOR CLASS statement in the SQL standard.
See Also

CREATE OPERATOR CLASS, DROP OPERATOR CLASS, ALTER OPERATOR FAMILY
ALTER OPERATOR FAMILY

Name
ALTER OPERATOR FAMILY — change the definition of an operator family

 Synopsis
ALTER OPERATOR FAMILY name USING index_method ADD
   { OPERATOR strategy_number operator_name ( op_type, op_type )
     [ FOR SEARCH | FOR ORDER BY sort_family_name ]
     | FUNCTION support_number [ ( op_type [ , op_type ] ) ]
     function_name ( argument_type [, ...] )
   } [, ... ]

ALTER OPERATOR FAMILY name USING index_method DROP
   { OPERATOR strategy_number ( op_type [ , op_type ] )
     | FUNCTION support_number ( op_type [ , op_type ] )
   } [, ... ]

ALTER OPERATOR FAMILY name USING index_method
   RENAME TO new_name

ALTER OPERATOR FAMILY name USING index_method
   OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

ALTER OPERATOR FAMILY name USING index_method
   SET SCHEMA new_schema

Description
ALTER OPERATOR FAMILY changes the definition of an operator family. You can add operators and support functions to the family, remove them from the family, or change the family’s name or owner.

When operators and support functions are added to a family with ALTER OPERATOR FAMILY, they are not part of any specific operator class within the family, but are just “loose” within the family. This indicates that these operators and functions are compatible with the family’s semantics, but are not required for correct functioning of any specific index. (Operators and functions that are so required should be declared as part of an operator class, instead; see CREATE OPERATOR CLASS.) PostgreSQL will allow loose members of a family to be dropped from the family at any time, but members of an operator class cannot be dropped without dropping the whole class and any indexes that depend on it. Typically, single-data-type operators and functions are part of operator classes because they are needed to support an index on that specific data type, while cross-data-type operators and functions are made loose members of the family.

You must be a superuser to use ALTER OPERATOR FAMILY. (This restriction is made because an erroneous operator family definition could confuse or even crash the server.)

ALTER OPERATOR FAMILY does not presently check whether the operator family definition includes all the operators and functions required by the index method, nor whether the operators and functions form a self-consistent set. It is the user’s responsibility to define a valid operator family.

Refer to Section 36.14 for further information.
Parameters

name
The name (optionally schema-qualified) of an existing operator family.

index_method
The name of the index method this operator family is for.

strategy_number
The index method’s strategy number for an operator associated with the operator family.

operator_name
The name (optionally schema-qualified) of an operator associated with the operator family.

op_type
In an OPERATOR clause, the operand data type(s) of the operator, or NONE to signify a left-unary or right-unary operator. Unlike the comparable syntax in CREATE OPERATOR CLASS, the operand data types must always be specified.

In an ADD FUNCTION clause, the operand data type(s) the function is intended to support, if different from the input data type(s) of the function. For B-tree comparison functions and hash functions it is not necessary to specify op_type since the function’s input data type(s) are always the correct ones to use. For B-tree sort support functions and all functions in GiST, SP-GiST and GIN operator classes, it is necessary to specify the operand data type(s) the function is to be used with.

In a DROP FUNCTION clause, the operand data type(s) the function is intended to support must be specified.

sort_family_name
The name (optionally schema-qualified) of an existing btree operator family that describes the sort ordering associated with an ordering operator.

If neither FOR SEARCH nor FOR ORDER BY is specified, FOR SEARCH is the default.

support_number
The index method’s support procedure number for a function associated with the operator family.

function_name
The name (optionally schema-qualified) of a function that is an index method support procedure for the operator family.

argument_type
The parameter data type(s) of the function.

new_name
The new name of the operator family.

new_owner
The new owner of the operator family.

new_schema
The new schema for the operator family.

The OPERATOR and FUNCTION clauses can appear in any order.
**Notes**

Notice that the `DROP` syntax only specifies the “slot” in the operator family, by strategy or support number and input data type(s). The name of the operator or function occupying the slot is not mentioned. Also, for `DROP FUNCTION` the type(s) to specify are the input data type(s) the function is intended to support; for GiST, SP-GiST and GIN indexes this might have nothing to do with the actual input argument types of the function.

Because the index machinery does not check access permissions on functions before using them, including a function or operator in an operator family is tantamount to granting public execute permission on it. This is usually not an issue for the sorts of functions that are useful in an operator family.

The operators should not be defined by SQL functions. A SQL function is likely to be inlined into the calling query, which will prevent the optimizer from recognizing that the query matches an index.

Before PostgreSQL 8.4, the `OPERATOR` clause could include a `RECHECK` option. This is no longer supported because whether an index operator is “lossy” is now determined on-the-fly at run time. This allows efficient handling of cases where an operator might or might not be lossy.

**Examples**

The following example command adds cross-data-type operators and support functions to an operator family that already contains B-tree operator classes for data types `int4` and `int2`.

```
ALTER OPERATOR FAMILY integer_ops USING btree ADD

-- int4 vs int2
OPERATOR 1 < (int4, int2) ,
OPERATOR 2 <= (int4, int2) ,
OPERATOR 3 = (int4, int2) ,
OPERATOR 4 >= (int4, int2) ,
OPERATOR 5 > (int4, int2) ,
FUNCTION 1 btint42cmp(int4, int2) ,

-- int2 vs int4
OPERATOR 1 < (int2, int4) ,
OPERATOR 2 <= (int2, int4) ,
OPERATOR 3 = (int2, int4) ,
OPERATOR 4 >= (int2, int4) ,
OPERATOR 5 > (int2, int4) ,
FUNCTION 1 btint24cmp(int2, int4) ;
```

To remove these entries again:

```
ALTER OPERATOR FAMILY integer_ops USING btree DROP

-- int4 vs int2
OPERATOR 1 (int4, int2) ,
OPERATOR 2 (int4, int2) ,
OPERATOR 3 (int4, int2) ,
OPERATOR 4 (int4, int2) ,
OPERATOR 5 (int4, int2) ,
FUNCTION 1 (int4, int2) ,

-- int2 vs int4
```
ALTER OPERATOR FAMILY

OPERATOR 1 (int2, int4) ,
OPERATOR 2 (int2, int4) ,
OPERATOR 3 (int2, int4) ,
OPERATOR 4 (int2, int4) ,
OPERATOR 5 (int2, int4) ,
FUNCTION 1 (int2, int4) ;

Compatibility

There is no ALTER OPERATOR FAMILY statement in the SQL standard.

See Also

CREATE OPERATOR FAMILY, DROP OPERATOR FAMILY, CREATE OPERATOR CLASS, ALTER OPERATOR CLASS, DROP OPERATOR CLASS
ALTER POLICY

Name
ALTER POLICY — change the definition of a row level security policy

Synopsis
ALTER POLICY name ON table_name RENAME TO new_name

ALTER POLICY name ON table_name
    [ TO { role_name | PUBLIC | CURRENT_USER | SESSION_USER } [, ...] ]
    [ USING ( using_expression ) ]
    [ WITH CHECK ( check_expression ) ]

Description
ALTER POLICY changes the definition of an existing row-level security policy.
To use ALTER POLICY, you must own the table that the policy applies to.
In the second form of ALTER POLICY, the role list, using_expression, and check_expression
are replaced independently if specified. When one of those clauses is omitted, the corresponding part
of the policy is unchanged.

Parameters
name
    The name of an existing policy to alter.
table_name
    The name (optionally schema-qualified) of the table that the policy is on.
new_name
    The new name for the policy.
role_name
    The role(s) to which the policy applies. Multiple roles can be specified at one time. To apply the
    policy to all roles, use PUBLIC.
using_expression
    The USING expression for the policy. See CREATE POLICY for details.
check_expression
    The WITH CHECK expression for the policy. See CREATE POLICY for details.
Compatibility

ALTER POLICY is a PostgreSQL extension.

See Also

CREATE POLICY, DROP POLICY
ALTER ROLE

Name
ALTER ROLE — change a database role

Synopsis
ALTER ROLE role_specification [ WITH ] option [ ... ]

where option can be:

SUPERUSER | NOSUPERUSER
| CREATEDB | NOCREATEDB
| CREATEROLE | NOCREATEROLE
| INHERIT | NOINHERIT
| LOGIN | NOLOGIN
| REPLICA | NOREPlication
| BYPASSRLS | NOBYPASSRLS
| CONNECTION LIMIT conlimit
| [ ENCRYPTED | UNENCRYPTED ] PASSWORD 'password'
| VALID UNTIL 'timestamp'

ALTER ROLE name RENAME TO new_name

ALTER ROLE { role_specification | ALL } [ IN DATABASE database_name ] SET configuration_parameter { TO | = } { value | DEFAULT }
ALTER ROLE { role_specification | ALL } [ IN DATABASE database_name ] RESET configuration_parameter
ALTER ROLE { role_specification | ALL } [ IN DATABASE database_name ] RESET ALL

where role_specification can be:

role_name
| CURRENT_USER
| SESSION_USER

Description
ALTER ROLE changes the attributes of a PostgreSQL role.

The first variant of this command listed in the synopsis can change many of the role attributes that can be specified in CREATE ROLE. (All the possible attributes are covered, except that there are no options for adding or removing memberships; use GRANT and REVOKE for that.) Attributes not mentioned in the command retain their previous settings. Database superusers can change any of these settings for any role. Roles having CREATEROLE privilege can change any of these settings, but only for non-superuser and non-replication roles. Ordinary roles can only change their own password.

The second variant changes the name of the role. Database superusers can rename any role. Roles having CREATEROLE privilege can rename non-superuser roles. The current session user cannot be renamed. (Connect as a different user if you need to do that.) Because MD5-encrypted passwords use the role name as cryptographic salt, renaming a role clears its password if the password is MD5-encrypted.
The remaining variants change a role’s session default for a configuration variable, either for all databases or, when the \texttt{IN DATABASE} clause is specified, only for sessions in the named database. If \texttt{ALL} is specified instead of a role name, this changes the setting for all roles. Using \texttt{ALL with IN DATABASE} is effectively the same as using the command \texttt{ALTER DATABASE ... SET ...}.

Whenever the role subsequently starts a new session, the specified value becomes the session default, overriding whatever setting is present in \texttt{postgresql.conf} or has been received from the \texttt{postgres} command line. This only happens at login time; executing \texttt{SET ROLE} or \texttt{SET SESSION AUTHORIZATION} does not cause new configuration values to be set. Settings set for all databases are overridden by database-specific settings attached to a role. Settings for specific databases or specific roles override settings for all roles.

Superusers can change anyone’s session defaults. Roles having \texttt{CREATEROLE} privilege can change defaults for non-superuser roles. Ordinary roles can only set defaults for themselves. Certain configuration variables cannot be set this way, or can only be set if a superuser issues the command. Only superusers can change a setting for all roles in all databases.

### Parameters

\texttt{name}

The name of the role whose attributes are to be altered.

\texttt{CURRENT_USER}

Alter the current user instead of an explicitly identified role.

\texttt{SESSION_USER}

Alter the current session user instead of an explicitly identified role.

\texttt{SUPERUSER}
\texttt{NOSUPERUSER}
\texttt{CREATEDB}
\texttt{NOCREATEDB}
\texttt{CREATEROLE}
\texttt{NOCREATEROLE}
\texttt{INHERIT}
\texttt{NOINHERIT}
\texttt{LOGIN}
\texttt{NOLOGIN}
\texttt{REPLICATION}
\texttt{NOREPLICATION}
\texttt{BYPASSRLS}
\texttt{NOBYPASSRLS}
\texttt{CONNECTION LIMIT connlimit}
\texttt{PASSWORD password}
\texttt{ENCRYPTED}
\texttt{UNENCRYPTED}
\texttt{VALID UNTIL ’timestamp’}

These clauses alter attributes originally set by \texttt{CREATE ROLE}. For more information, see the \texttt{CREATE ROLE} reference page.

\texttt{new_name}

The new name of the role.
ALTER ROLE  

database_name  
The name of the database the configuration variable should be set in.  

configuration_parameter  
value  
Set this role's session default for the specified configuration parameter to the given value. If value is DEFAULT or, equivalently, RESET is used, the role-specific variable setting is removed, so the role will inherit the system-wide default setting in new sessions. Use RESET ALL to clear all role-specific settings. SET FROM CURRENT saves the session’s current value of the parameter as the role-specific value. If IN DATABASE is specified, the configuration parameter is set or removed for the given role and database only.  

Role-specific variable settings take effect only at login; SET ROLE and SET SESSION AUTHORIZATION do not process role-specific variable settings.  

See SET and Chapter 19 for more information about allowed parameter names and values.  

Notes  
Use CREATE ROLE to add new roles, and DROP ROLE to remove a role.  

ALTER ROLE cannot change a role’s memberships. Use GRANT and REVOKE to do that.  

Caution must be exercised when specifying an unencrypted password with this command. The password will be transmitted to the server in cleartext, and it might also be logged in the client’s command history or the server log. psql contains a command \password that can be used to change a role’s password without exposing the cleartext password.  

It is also possible to tie a session default to a specific database rather than to a role; see ALTER DATABASE. If there is a conflict, database-role-specific settings override role-specific ones, which in turn override database-specific ones.  

Examples  
Change a role’s password:  

ALTER ROLE davide WITH PASSWORD ‘hu8jmn3’;  

Remove a role’s password:  

ALTER ROLE davide WITH PASSWORD NULL;  

Change a password expiration date, specifying that the password should expire at midday on 4th May 2015 using the time zone which is one hour ahead of UTC:  

ALTER ROLE chris VALID UNTIL ‘May 4 12:00:00 2015 +1’;  

Make a password valid forever:  

ALTER ROLE fred VALID UNTIL ‘infinity’;
Give a role the ability to create other roles and new databases:

```
ALTER ROLE miriam CREATEROLE CREATEDB;
```

Give a role a non-default setting of the maintenance_work_mem parameter:

```
ALTER ROLE worker_bee SET maintenance_work_mem = 100000;
```

Give a role a non-default, database-specific setting of the client_min_messages parameter:

```
ALTER ROLE fred IN DATABASE devel SET client_min_messages = DEBUG;
```

**Compatibility**

The `ALTER ROLE` statement is a PostgreSQL extension.

**See Also**

`CREATE ROLE`, `DROP ROLE`, `ALTER DATABASE`, `SET`
ALTER RULE

Name
ALTER RULE — change the definition of a rule

Synopsis
ALTER RULE name ON table_name RENAME TO new_name

Description
ALTER RULE changes properties of an existing rule. Currently, the only available action is to change the rule’s name.
To use ALTER RULE, you must own the table or view that the rule applies to.

Parameters

name
   The name of an existing rule to alter.

table_name
   The name (optionally schema-qualified) of the table or view that the rule applies to.

new_name
   The new name for the rule.

Examples
To rename an existing rule:
ALTER RULE notify_all ON emp RENAME TO notify_me;

Compatibility
ALTER RULE is a PostgreSQL language extension, as is the entire query rewrite system.

See Also
CREATE RULE, DROP RULE
ALTER SCHEMA

Name

ALTER SCHEMA — change the definition of a schema

Synopsis

ALTER SCHEMA name RENAME TO new_name
ALTER SCHEMA name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }

Description

ALTER SCHEMA changes the definition of a schema.

You must own the schema to use ALTER SCHEMA. To rename a schema you must also have the CREATE privilege for the database. To alter the owner, you must also be a direct or indirect member of the new owning role, and you must have the CREATE privilege for the database. (Note that superusers have all these privileges automatically.)

Parameters

name

The name of an existing schema.

new_name

The new name of the schema. The new name cannot begin with pg_, as such names are reserved for system schemas.

new_owner

The new owner of the schema.

Compatibility

There is no ALTER SCHEMA statement in the SQL standard.

See Also

CREATE SCHEMA, DROP SCHEMA
ALTER SEQUENCE

Name
ALTER SEQUENCE — change the definition of a sequence generator

Synopsis

ALTER SEQUENCE [ IF EXISTS ] name [ INCREMENT [ BY ] increment ]
[ MINVALUE minvalue | NO MINVALUE ] [ MAXVALUE maxvalue | NO MAXVALUE ]
[ START [ WITH ] start ]
[ RESTART [ [ WITH ] restart ] ]
[ CACHE cache ] [ [ NO ] CYCLE ]
[ OWNED BY { table_name.column_name | NONE } ]
ALTER SEQUENCE [ IF EXISTS ] name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER SEQUENCE [ IF EXISTS ] name RENAME TO new_name
ALTER SEQUENCE [ IF EXISTS ] name SET SCHEMA new_schema

Description

ALTER SEQUENCE changes the parameters of an existing sequence generator. Any parameters not
specifically set in the ALTER SEQUENCE command retain their prior settings.

You must own the sequence to use ALTER SEQUENCE. To change a sequence’s schema, you must also
have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect
member of the new owning role, and that role must have CREATE privilege on the sequence’s schema.
(These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping
and recreating the sequence. However, a superuser can alter ownership of any sequence anyway.)

Parameters

name

The name (optionally schema-qualified) of a sequence to be altered.

IF EXISTS

Do not throw an error if the sequence does not exist. A notice is issued in this case.

increment

The clause INCREMENT BY increment is optional. A positive value will make an ascending
sequence, a negative one a descending sequence. If unspecified, the old increment value will be
maintained.

minvalue

NO MINVALUE

The optional clause MINVALUE minvalue determines the minimum value a sequence can gen-
erate. If NO MINVALUE is specified, the defaults of 1 and -2\(^{31}\)-1 for ascending and descending
sequences, respectively, will be used. If neither option is specified, the current minimum value
will be maintained.
ALTER SEQUENCE

maxvalue
NO MAXVALUE

The optional clause MAXVALUE maxvalue determines the maximum value for the sequence. If NO MAXVALUE is specified, the defaults are $2^{63}$-1 and -1 for ascending and descending sequences, respectively, will be used. If neither option is specified, the current maximum value will be maintained.

start

The optional clause START WITH start changes the recorded start value of the sequence. This has no effect on the current sequence value; it simply sets the value that future ALTER SEQUENCE commands will use.

restart

The optional clause RESTART [ WITH restart ] changes the current value of the sequence. This is equivalent to calling the setval function with is_called = false: the specified value will be returned by the next call of nextval. Writing RESTART with no restart value is equivalent to supplying the start value that was recorded by CREATE SEQUENCE or last set by ALTER SEQUENCE START WITH.

cache

The clause CACHE cache enables sequence numbers to be preallocated and stored in memory for faster access. The minimum value is 1 (only one value can be generated at a time, i.e., no cache). If unspecified, the old cache value will be maintained.

CYCLE

The optional CYCLE key word can be used to enable the sequence to wrap around when the maxvalue or minvalue has been reached by an ascending or descending sequence respectively. If the limit is reached, the next number generated will be the minvalue or maxvalue, respectively.

NO CYCLE

If the optional NO CYCLE key word is specified, any calls to nextval after the sequence has reached its maximum value will return an error. If neither CYCLE or NO CYCLE are specified, the old cycle behavior will be maintained.

OWNED BY table_name.column_name
OWNED BY NONE

The OWNED BY option causes the sequence to be associated with a specific table column, such that if that column (or its whole table) is dropped, the sequence will be automatically dropped as well. If specified, this association replaces any previously specified association for the sequence. The specified table must have the same owner and be in the same schema as the sequence. Specifying OWNED BY NONE removes any existing association, making the sequence “free-standing”.

new_owner

The user name of the new owner of the sequence.

new_name

The new name for the sequence.

new_schema

The new schema for the sequence.
### Notes

To avoid blocking of concurrent transactions that obtain numbers from the same sequence, `ALTER SEQUENCE`’s effects on the sequence generation parameters are never rolled back; those changes take effect immediately and are not reversible. However, the `OWNED BY`, `OWNER TO`, `RENAME TO`, and `SET SCHEMA` clauses cause ordinary catalog updates that can be rolled back.

`ALTER SEQUENCE` will not immediately affect `nextval` results in backends, other than the current one, that have preallocated (cached) sequence values. They will use up all cached values prior to noticing the changed sequence generation parameters. The current backend will be affected immediately.

`ALTER SEQUENCE` does not affect the `currval` status for the sequence. (Before PostgreSQL 8.3, it sometimes did.)

For historical reasons, `ALTER TABLE` can be used with sequences too; but the only variants of `ALTER TABLE` that are allowed with sequences are equivalent to the forms shown above.

### Examples

Restart a sequence called `serial`, at 105:

```
ALTER SEQUENCE serial RESTART WITH 105;
```

### Compatibility

`ALTER SEQUENCE` conforms to the SQL standard, except for the `START WITH`, `OWNED BY`, `OWNER TO`, `RENAME TO`, and `SET SCHEMA` clauses, which are PostgreSQL extensions.

### See Also

`CREATE SEQUENCE`, `DROP SEQUENCE`
ALTER SERVER

Name
ALTER SERVER — change the definition of a foreign server

Synopsis
ALTER SERVER name [ VERSION 'new_version' ]
   [ OPTIONS ([ ADD | SET | DROP ] option ['value'] [, ... ]) ]
ALTER SERVER name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER SERVER name RENAME TO new_name

Description
ALTER SERVER changes the definition of a foreign server. The first form changes the server version string or the generic options of the server (at least one clause is required). The second form changes the owner of the server.

To alter the server you must be the owner of the server. Additionally to alter the owner, you must own the server and also be a direct or indirect member of the new owning role, and you must have USAGE privilege on the server’s foreign-data wrapper. (Note that superusers satisfy all these criteria automatically.)

Parameters
name
   The name of an existing server.
new_version
   New server version.
OPTIONS ( [ ADD | SET | DROP ] option ['value'] [, ... ] )
   Change options for the server. ADD, SET, and DROP specify the action to be performed. ADD is assumed if no operation is explicitly specified. Option names must be unique; names and values are also validated using the server’s foreign-data wrapper library.
new_owner
   The user name of the new owner of the foreign server.
new_name
   The new name for the foreign server.

Examples
Alter server foo, add connection options:

ALTER SERVER foo OPTIONS (host 'foo', dbname 'foodb');
Alter server foo, change version, change host option:

```
ALTER SERVER foo VERSION '8.4' OPTIONS (SET host 'baz');
```

**Compatibility**

`ALTER SERVER` conforms to ISO/IEC 9075-9 (SQL/MED). The `OWNER TO` and `RENAME` forms are PostgreSQL extensions.

**See Also**

`CREATE SERVER`, `DROP SERVER`
ALTER SYSTEM

Name
ALTER SYSTEM — change a server configuration parameter

Synopsis
ALTER SYSTEM SET configuration_parameter { TO | = } { value | 'value' | DEFAULT }
ALTER SYSTEM RESET configuration_parameter
ALTER SYSTEM RESET ALL

Description
ALTER SYSTEM is used for changing server configuration parameters across the entire database cluster. It can be more convenient than the traditional method of manually editing the postgresql.conf file. ALTER SYSTEM writes the given parameter setting to the postgresql.auto.conf file, which is read in addition to postgresql.conf. Setting a parameter to DEFAULT, or using the RESET variant, removes that configuration entry from the postgresql.auto.conf file. Use RESET ALL to remove all such configuration entries.

Values set with ALTER SYSTEM will be effective after the next server configuration reload, or after the next server restart in the case of parameters that can only be changed at server start. A server configuration reload can be commanded by calling the SQL function pg_reload_conf(), running pg_ctl reload, or sending a SIGHUP signal to the main server process.

Only superusers can use ALTER SYSTEM. Also, since this command acts directly on the file system and cannot be rolled back, it is not allowed inside a transaction block or function.

Parameters

configuration_parameter

Name of a settable configuration parameter. Available parameters are documented in Chapter 19.

value

New value of the parameter. Values can be specified as string constants, identifiers, numbers, or comma-separated lists of these, as appropriate for the particular parameter. DEFAULT can be written to specify removing the parameter and its value from postgresql.auto.conf.

Notes

This command can’t be used to set data_directory, nor parameters that are not allowed in postgresql.conf (e.g., preset options).

See Section 19.1 for other ways to set the parameters.
ALTER SYSTEM

Examples

Set the `wal_level`:

```
ALTER SYSTEM SET wal_level = replica;
```

Undo that, restoring whatever setting was effective in `postgresql.conf`:

```
ALTER SYSTEM RESET wal_level;
```

Compatibility

The `ALTER SYSTEM` statement is a PostgreSQL extension.

See Also

`SET`, `SHOW`
**ALTER TABLE**

**Name**

ALTER TABLE — change the definition of a table

**Synopsis**

ALTER TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
   action [ , ... ]
ALTER TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
   RENAME [ COLUMN ] column_name TO new_column_name
ALTER TABLE [ IF EXISTS ] [ ONLY ] name [ * ]
   RENAME CONSTRAINT constraint_name TO new_constraint_name
ALTER TABLE [ IF EXISTS ] name
   RENAME TO new_name
ALTER TABLE [ IF EXISTS ] name
   SET SCHEMA new_schema
ALTER TABLE ALL IN TABLESPACE name [ OWNED BY role_name [ , ... ] ]
   SET TABLESPACE new_tablespace [ NOWAIT ]

where action is one of:

ADD [ COLUMN ] [ IF NOT EXISTS ] column_name data_type [ COLLATE collation ] [ column_constraint ]
DROP [ COLUMN ] [ IF EXISTS ] column_name [ RESTRICT | CASCADE ]
ALTER [ COLUMN ] column_name [ SET DATA ] TYPE data_type [ COLLATE collation ] [ USING expression ]
ALTER [ COLUMN ] column_name SET DEFAULT expression
ALTER [ COLUMN ] column_name DROP DEFAULT
ALTER [ COLUMN ] column_name ( SET | DROP ) NOT NULL
ALTER [ COLUMN ] column_name SET STATISTICS integer
ALTER [ COLUMN ] column_name SET ( attribute_option = value [ , ... ] )
ALTER [ COLUMN ] column_name RESET ( attribute_option [ , ... ] )
ALTER [ COLUMN ] column_name SET STORAGE { PLAIN | EXTERNAL | EXTENDED | MAIN }
ADD table_constraint [ NOT VALID ]
ADD table_constraint_using_index
ALTER CONSTRAINT constraint_name [ DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED ]
VALIDATE CONSTRAINT constraint_name
DROP CONSTRAINT [ IF EXISTS ] constraint_name [ RESTRICT | CASCADE ]
DISABLE TRIGGER [ trigger_name | ALL | USER ]
ENABLE TRIGGER [ trigger_name | ALL | USER ]
ENABLE REPLICA TRIGGER trigger_name
ENABLE ALWAYS TRIGGER trigger_name
DISABLE RULE rewrite_rule_name
ENABLE RULE rewrite_rule_name
ENABLE REPLICA RULE rewrite_rule_name
ENABLE ALWAYS RULE rewrite_rule_name
DISABLE ROW LEVEL SECURITY
ENABLE ROW LEVEL SECURITY
FORCE ROW LEVEL SECURITY
NO FORCE ROW LEVEL SECURITY
CLUSTER ON index_name
SET WITHOUT CLUSTER
SET WITH OIDS
SET WITHOUT OIDS
SET TABLESPACE new_tablespace
ALTERTABLE

SET { LOGGED | UNLOGGED }
SET { storage_parameter = value [, ... ] }
RESET { storage_parameter [, ... ] }
INHERIT parent_table
NO INHERIT parent_table
OF type_name
NOT OF
OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
REPLICA IDENTITY { DEFAULT | USING INDEX index_name | FULL | NOTHING }

and table_constraint_using_index is:

[ CONSTRAINT constraint_name ]
{ UNIQUE | PRIMARY KEY } USING INDEX index_name
[ DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ]

Description

ALTERTABLE changes the definition of an existing table. There are several subforms described below. Note that the lock level required may differ for each subform. An ACCESS EXCLUSIVE lock is held unless explicitly noted. When multiple subcommands are listed, the lock held will be the strictest one required from any subcommand.

ADD COLUMN [ IF NOT EXISTS ]

This form adds a new column to the table, using the same syntax as CREATE TABLE. If IF NOT EXISTS is specified and a column already exists with this name, no error is thrown.

DROP COLUMN [ IF EXISTS ]

This form drops a column from a table. Indexes and table constraints involving the column will be automatically dropped as well. You will need to say CASCADE if anything outside the table depends on the column, for example, foreign key references or views. If IF EXISTS is specified and the column does not exist, no error is thrown. In this case a notice is issued instead.

SET DATA TYPE

This form changes the type of a column of a table. Indexes and simple table constraints involving the column will be automatically converted to use the new column type by reparsing the originally supplied expression. The optional COLLATE clause specifies a collation for the new column; if omitted, the collation is the default for the new column type. The optional USING clause specifies how to compute the new column value from the old; if omitted, the default conversion is the same as an assignment cast from old data type to new. A USING clause must be provided if there is no implicit or assignment cast from old to new type.

SET/DROP DEFAULT

These forms set or remove the default value for a column. Default values only apply in subsequent INSERT or UPDATE commands; they do not cause rows already in the table to change.

SET/DROP NOT NULL

These forms change whether a column is marked to allow null values or to reject null values. You can only use SET NOT NULL when the column contains no null values.
SET STATISTICS

This form sets the per-column statistics-gathering target for subsequent ANALYZE operations. The target can be set in the range 0 to 10000; alternatively, set it to -1 to revert to using the system default statistics target (default_statistics_target). For more information on the use of statistics by the PostgreSQL query planner, refer to Section 14.2.

SET STATISTICS acquires a SHARE UPDATE EXCLUSIVE lock.

SET (attribute_option = value [, ... ])
RESET (attribute_option [, ... ])

This form sets or resets per-attribute options. Currently, the only defined per-attribute options are n_distinct and n_distinct_inherited, which override the number-of-distinct-values estimates made by subsequent ANALYZE operations. n_distinct affects the statistics for the table itself, while n_distinct_inherited affects the statistics gathered for the table plus its inheritance children. When set to a positive value, ANALYZE will assume that the column contains exactly the specified number of distinct nonnull values. When set to a negative value, which must be greater than or equal to -1, ANALYZE will assume that the number of distinct nonnull values in the column is linear in the size of the table; the exact count is to be computed by multiplying the estimated table size by the absolute value of the given number. For example, a value of -1 implies that all values in the column are distinct, while a value of -0.5 implies that each value appears twice on the average. This can be useful when the size of the table changes over time, since the multiplication by the number of rows in the table is not performed until query planning time. Specify a value of 0 to revert to estimating the number of distinct values normally. For more information on the use of statistics by the PostgreSQL query planner, refer to Section 14.2.

Changing per-attribute options acquires a SHARE UPDATE EXCLUSIVE lock.

SET STORAGE

This form sets the storage mode for a column. This controls whether this column is held inline or in a secondary TOAST table, and whether the data should be compressed or not. PLAIN must be used for fixed-length values such as integer and is inline, uncompressed. MAIN is for inline, compressible data. EXTERNAL is for external, uncompressed data, and EXTENDED is for external, compressed data. EXTENDED is the default for most data types that support non-PLAIN storage. Use of EXTERNAL will make substring operations on very large text and bytea values run faster, at the penalty of increased storage space. Note that SET STORAGE doesn’t itself change anything in the table, it just sets the strategy to be pursued during future table updates. See Section 65.2 for more information.

ADD table_constraint [ NOT VALID ]

This form adds a new constraint to a table using the same syntax as CREATE TABLE, plus the option NOT VALID, which is currently only allowed for foreign key and CHECK constraints.

Normally, this form will cause a scan of the table to verify that all existing rows in the table satisfy the new constraint. But if the NOT VALID option is used, this potentially-lengthy scan is skipped. The constraint will still be enforced against subsequent inserts or updates (that is, they’ll fail unless there is a matching row in the referenced table, in the case of foreign keys, or they’ll fail unless the new row matches the specified check condition). But the database will not assume that the constraint holds for all rows in the table, until it is validated by using the VALIDATE CONSTRAINT option. See Notes below for more information about using the NOT VALID option.

ADD table_constraint_using_index

This form adds a new PRIMARY KEY or UNIQUE constraint to a table based on an existing unique index. All the columns of the index will be included in the constraint.
ALTER TABLE

The index cannot have expression columns nor be a partial index. Also, it must be a b-tree index
with default sort ordering. These restrictions ensure that the index is equivalent to one that would
be built by a regular \texttt{ADD PRIMARY KEY} or \texttt{ADD UNIQUE} command.

If \texttt{PRIMARY KEY} is specified, and the index’s columns are not already marked \texttt{NOT NULL}, then
this command will attempt to do \texttt{ALTER COLUMN SET NOT NULL} against each such column.
That requires a full table scan to verify the column(s) contain no nulls. In all other cases, this is
a fast operation.

If a constraint name is provided then the index will be renamed to match the constraint name.
Otherwise the constraint will be named the same as the index.

After this command is executed, the index is “owned” by the constraint, in the same way as if the
index had been built by a regular \texttt{ADD PRIMARY KEY} or \texttt{ADD UNIQUE} command. In particular,
dropping the constraint will make the index disappear too.

\textbf{Note:} Adding a constraint using an existing index can be helpful in situations where a new
constraint needs to be added without blocking table updates for a long time. To do that, create
the index using \texttt{CREATE INDEX CONCURRENTLY}, and then install it as an official constraint
using this syntax. See the example below.

\textbf{ALTER CONSTRAINT}

This form alters the attributes of a constraint that was previously created. Currently only foreign
key constraints may be altered.

\textbf{VALIDATE CONSTRAINT}

This form validates a foreign key or check constraint that was previously created as \texttt{NOT VALID},
by scanning the table to ensure there are no rows for which the constraint is not satisfied. Nothing
happens if the constraint is already marked valid. (See \textit{Notes} below for an explanation of the
usefulness of this command.)

\textbf{DROP CONSTRAINT [ IF EXISTS ]}

This form drops the specified constraint on a table. If \texttt{IF EXISTS} is specified and the constraint
does not exist, no error is thrown. In this case a notice is issued instead.

\textbf{DISABLE/ENABLE [ REPLICA | ALWAYS ] TRIGGER}

These forms configure the firing of trigger(s) belonging to the table. A disabled trigger is still
known to the system, but is not executed when its triggering event occurs. For a deferred trigger,
the enable status is checked when the event occurs, not when the trigger function is actually
executed. One can disable or enable a single trigger specified by name, or all triggers on the
table, or only user triggers (this option excludes internally generated constraint triggers such as
those that are used to implement foreign key constraints or deferrable uniqueness and exclusion
constraints). Disabling or enabling internally generated constraint triggers requires superuser
privileges; it should be done with caution since of course the integrity of the constraint cannot
be guaranteed if the triggers are not executed. The trigger firing mechanism is also affected by
the configuration variable \texttt{session_replication_role}. Simply enabled triggers will fire when the
replication role is “origin” (the default) or “local”. Triggers configured as \texttt{ENABLE REPLICA}
will only fire if the session is in “replica” mode, and triggers configured as \texttt{ENABLE ALWAYS}
will fire regardless of the current replication mode.

This command acquires a \texttt{SHARE ROW EXCLUSIVE} lock.
ALTER TABLE

DISABLE/ENABLE [ REPLICA | ALWAYS ] RULE

These forms configure the firing of rewrite rules belonging to the table. A disabled rule is still known to the system, but is not applied during query rewriting. The semantics are as for disabled/enabled triggers. This configuration is ignored for ON SELECT rules, which are always applied in order to keep views working even if the current session is in a non-default replication role.

DISABLE/ENABLE ROW LEVEL SECURITY

These forms control the application of row security policies belonging to the table. If enabled and no policies exist for the table, then a default-deny policy is applied. Note that policies can exist for a table even if row level security is disabled - in this case, the policies will NOT be applied and the policies will be ignored. See also CREATE POLICY.

NO FORCE/FORCE ROW LEVEL SECURITY

These forms control the application of row security policies belonging to the table when the user is the table owner. If enabled, row level security policies will be applied when the user is the table owner. If disabled (the default) then row level security will not be applied when the user is the table owner. See also CREATE POLICY.

CLUSTER ON

This form selects the default index for future CLUSTER operations. It does not actually re-cluster the table.

Changing cluster options acquires a SHARE UPDATE EXCLUSIVE lock.

SET WITHOUT CLUSTER

This form removes the most recently used CLUSTER index specification from the table. This affects future cluster operations that don’t specify an index.

Changing cluster options acquires a SHARE UPDATE EXCLUSIVE lock.

SET WITH OIDS

This form adds an oid system column to the table (see Section 5.4). It does nothing if the table already has OIDs.

Note that this is not equivalent to ADD COLUMN oid oid; that would add a normal column that happened to be named oid, not a system column.

SET WITHOUT OIDS

This form removes the oid system column from the table. This is exactly equivalent to DROP COLUMN oid RESTRICT, except that it will not complain if there is already no oid column.

SET TABLESPACE

This form changes the table’s tablespace to the specified tablespace and moves the data file(s) associated with the table to the new tablespace. Indexes on the table, if any, are not moved; but they can be moved separately with additional SET TABLESPACE commands. All tables in the current database in a tablespace can be moved by using the ALL IN TABLESPACE form, which will lock all tables to be moved first and then move each one. This form also supports OWNED BY, which will only move tables owned by the roles specified. If the NOWAIT option is specified then the command will fail if it is unable to acquire all of the locks required immediately. Note that system catalogs are not moved by this command, use ALTER DATABASE or explicit ALTER TABLE invocations instead if desired. The information_schema relations are not considered part of the system catalogs and will be moved. See also CREATE TABLESPACE.
ALTER TABLE

SET \{ LOGGED | UNLOGGED \}

This form changes the table from unlogged to logged or vice-versa (see UNLOGGED). It cannot be applied to a temporary table.

SET \{ storage_parameter = value [, ... ] \}

This form changes one or more storage parameters for the table. See Storage Parameters for details on the available parameters. Note that the table contents will not be modified immediately by this command; depending on the parameter you might need to rewrite the table to get the desired effects. That can be done with VACUUM FULL, CLUSTER or one of the forms of ALTER TABLE that forces a table rewrite.

Changing fillfactor and autovacuum storage parameters acquires a SHARE UPDATE EXCLUSIVE lock.

Note: While CREATE TABLE allows OIDS to be specified in the WITH \{ storage_parameter \} syntax, ALTER TABLE does not treat OIDS as a storage parameter. Instead use the SET WITH OIDS and SET WITHOUT OIDS forms to change OID status.

RESET \{ storage_parameter [, ... ] \}

This form resets one or more storage parameters to their defaults. As with SET, a table rewrite might be needed to update the table entirely.

INHERIT parent_table

This form adds the target table as a new child of the specified parent table. Subsequently, queries against the parent will include records of the target table. To be added as a child, the target table must already contain all the same columns as the parent (it could have additional columns, too). The columns must have matching data types, and if they have NOT NULL constraints in the parent then they must also have NOT NULL constraints in the child.

There must also be matching child-table constraints for all CHECK constraints of the parent, except those marked non-inheritable (that is, created with ALTER TABLE ... ADD CONSTRAINT ... NO INHERIT) in the parent, which are ignored; all child-table constraints matched must not be marked non-inheritable. Currently UNIQUE, PRIMARY KEY, and FOREIGN KEY constraints are not considered, but this might change in the future.

NO INHERIT parent_table

This form removes the target table from the list of children of the specified parent table. Queries against the parent table will no longer include records drawn from the target table.

OF type_name

This form links the table to a composite type as though CREATE TABLE OF had formed it. The table’s list of column names and types must precisely match that of the composite type; the presence of an oid system column is permitted to differ. The table must not inherit from any other table. These restrictions ensure that CREATE TABLE OF would permit an equivalent table definition.

NOT OF

This form dissociates a typed table from its type.

OWNER

This form changes the owner of the table, sequence, view, materialized view, or foreign table to the specified user.
ALTER TABLE

REPLICA IDENTITY

This form changes the information which is written to the write-ahead log to identify rows which are updated or deleted. This option has no effect except when logical replication is in use. DEFAULT (the default for non-system tables) records the old values of the columns of the primary key, if any. USING INDEX records the old values of the columns covered by the named index, which must be unique, not partial, not deferrable, and include only columns marked NOT NULL. FULL records the old values of all columns in the row. NOTHING records no information about the old row. (This is the default for system tables.) In all cases, no old values are logged unless at least one of the columns that would be logged differs between the old and new versions of the row.

RENAME

The RENAME forms change the name of a table (or an index, sequence, view, materialized view, or foreign table), the name of an individual column in a table, or the name of a constraint of the table. There is no effect on the stored data.

SET SCHEMA

This form moves the table into another schema. Associated indexes, constraints, and sequences owned by table columns are moved as well.

All the forms of ALTER TABLE that act on a single table, except RENAME, and SET SCHEMA can be combined into a list of multiple alterations to applied together. For example, it is possible to add several columns and/or alter the type of several columns in a single command. This is particularly useful with large tables, since only one pass over the table need be made.

You must own the table to use ALTER TABLE. To change the schema or tablespace of a table, you must also have CREATE privilege on the new schema or tablespace. To add the table as a new child of a parent table, you must own the parent table as well. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the table’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the table. However, a superuser can alter ownership of any table anyway.) To add a column or alter a column type or use the OF clause, you must also have USAGE privilege on the data type.

Parameters

IF EXISTS

Do not throw an error if the table does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing table to alter. If ONLY is specified before the table name, only that table is altered. If ONLY is not specified, the table and all its descendant tables (if any) are altered. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

column_name

Name of a new or existing column.

new_column_name

New name for an existing column.
new_name
New name for the table.
data_type
Data type of the new column, or new data type for an existing column.
table_constraint
New table constraint for the table.
constraint_name
Name of a new or existing constraint.
CASCADE
Automatically drop objects that depend on the dropped column or constraint (for example, views referencing the column), and in turn all objects that depend on those objects (see Section 5.13).
RESTRICT
Refuse to drop the column or constraint if there are any dependent objects. This is the default behavior.
trigger_name
Name of a single trigger to disable or enable.
ALL
Disable or enable all triggers belonging to the table. (This requires superuser privilege if any of the triggers are internally generated constraint triggers such as those that are used to implement foreign key constraints or deferrable uniqueness and exclusion constraints.)
USER
Disable or enable all triggers belonging to the table except for internally generated constraint triggers such as those that are used to implement foreign key constraints or deferrable uniqueness and exclusion constraints.
index_name
The name of an existing index.
storage_parameter
The name of a table storage parameter.
value
The new value for a table storage parameter. This might be a number or a word depending on the parameter.
parent_table
A parent table to associate or de-associate with this table.
new_owner
The user name of the new owner of the table.
new_tablespace
The name of the tablespace to which the table will be moved.
The name of the schema to which the table will be moved.

Notes

The key word COLUMN is noise and can be omitted.

When a column is added with ADD COLUMN, all existing rows in the table are initialized with the column’s default value (NULL if no DEFAULT clause is specified). If there is no DEFAULT clause, this is merely a metadata change and does not require any immediate update of the table’s data; the added NULL values are supplied on readout, instead.

Adding a column with a DEFAULT clause or changing the type of an existing column will require the entire table and its indexes to be rewritten. As an exception when changing the type of an existing column, if the USING clause does not change the column contents and the old type is either binary coercible to the new type or an unconstrained domain over the new type, a table rewrite is not needed; but any indexes on the affected columns must still be rebuilt. Adding or removing a system oid column also requires rewriting the entire table. Table and/or index rebuilds may take a significant amount of time for a large table; and will temporarily require as much as double the disk space.

Adding a CHECK or NOT NULL constraint requires scanning the table to verify that existing rows meet the constraint, but does not require a table rewrite.

The main reason for providing the option to specify multiple changes in a single ALTER TABLE is that multiple table scans or rewrites can thereby be combined into a single pass over the table.

Scanning a large table to verify a new foreign key or check constraint can take a long time, and other updates to the table are locked out until the ALTER TABLE ADD CONSTRAINT command is committed. The main purpose of the NOT VALID constraint option is to reduce the impact of adding a constraint on concurrent updates. With NOT VALID, the ADD CONSTRAINT command does not scan the table and can be committed immediately. After that, a VALIDATE CONSTRAINT command can be issued to verify that existing rows satisfy the constraint. The validation step does not need to lock out concurrent updates, since it knows that other transactions will be enforcing the constraint for rows that they insert or update; only pre-existing rows need to be checked. Hence, validation acquires only a SHARE UPDATE EXCLUSIVE lock on the table being altered. (If the constraint is a foreign key then a ROW SHARE lock is also required on the table referenced by the constraint.) In addition to improving concurrency, it can be useful to use NOT VALID and VALIDATE CONSTRAINT in cases where the table is known to contain pre-existing violations. Once the constraint is in place, no new violations can be inserted, and the existing problems can be corrected at leisure until VALIDATE CONSTRAINT finally succeeds.

The DROP COLUMN form does not physically remove the column, but simply makes it invisible to SQL operations. Subsequent insert and update operations in the table will store a null value for the column. Thus, dropping a column is quick but it will not immediately reduce the on-disk size of your table, as the space occupied by the dropped column is not reclaimed. The space will be reclaimed over time as existing rows are updated. (These statements do not apply when dropping the system oid column; that is done with an immediate rewrite.)

To force immediate reclamation of space occupied by a dropped column, you can execute one of the forms of ALTER TABLE that performs a rewrite of the whole table. This results in reconstructing each row with the dropped column replaced by a null value.

The rewriting forms of ALTER TABLE are not MVCC-safe. After a table rewrite, the table will appear empty to concurrent transactions, if they are using a snapshot taken before the rewrite occurred. See Section 13.5 for more details.
The `USING` option of `SET DATA TYPE` can actually specify any expression involving the old values of the row; that is, it can refer to other columns as well as the one being converted. This allows very general conversions to be done with the `SET DATA TYPE` syntax. Because of this flexibility, the `USING` expression is not applied to the column’s default value (if any); the result might not be a constant expression as required for a default. This means that when there is no implicit or assignment cast from old to new type, `SET DATA TYPE` might fail to convert the default even though a `USING` clause is supplied. In such cases, drop the default with `DROP DEFAULT`, perform the `ALTER TYPE`, and then use `SET DEFAULT` to add a suitable new default. Similar considerations apply to indexes and constraints involving the column.

If a table has any descendant tables, it is not permitted to add, rename, or change the type of a column, or rename an inherited constraint in the parent table without doing the same to the descendants. That is, `ALTER TABLE ONLY` will be rejected. This ensures that the descendants always have columns matching the parent.

A recursive `DROP COLUMN` operation will remove a descendant table’s column only if the descendant does not inherit that column from any other parents and never had an independent definition of the column. A nonrecursive `DROP COLUMN` (i.e., `ALTER TABLE ONLY ... DROP COLUMN`) never removes any descendant columns, but instead marks them as independently defined rather than inherited.

The `TRIGGER`, `CLUSTER`, `OWNER`, and `TABLESPACE` actions never recurse to descendant tables; that is, they always act as though `ONLY` were specified. Adding a constraint recurses only for `CHECK` constraints that are not marked `NO INHERIT`.

Changing any part of a system catalog table is not permitted.

Refer to `CREATE TABLE` for a further description of valid parameters. Chapter 5 has further information on inheritance.

**Examples**

To add a column of type `varchar` to a table:

```sql
ALTER TABLE distributors ADD COLUMN address varchar(30);
```

To drop a column from a table:

```sql
ALTER TABLE distributors DROP COLUMN address RESTRICT;
```

To change the types of two existing columns in one operation:

```sql
ALTER TABLE distributors
    ALTER COLUMN address TYPE varchar(80),
    ALTER COLUMN name TYPE varchar(100);
```

To change an integer column containing Unix timestamps to `timestamp with time zone` via a `USING` clause:

```sql
ALTER TABLE foo
    ALTER COLUMN foo_timestamp SET DATA TYPE timestamp with time zone
    USING
        timestamp with time zone 'epoch' + foo_timestamp * interval '1 second';
```
ALTER TABLE foo
    ALTER COLUMN foo_timestamp DROP DEFAULT,
    ALTER COLUMN foo_timestamp TYPE timestamp with time zone
    USING
        timestamp with time zone 'epoch' + foo_timestamp * interval '1 second',
    ALTER COLUMN foo_timestamp SET DEFAULT now();

To rename an existing column:
ALTER TABLE distributors RENAME COLUMN address TO city;

To rename an existing table:
ALTER TABLE distributors RENAME TO suppliers;

To rename an existing constraint:
ALTER TABLE distributors RENAME CONSTRAINT zipchk TO zip_check;

To add a not-null constraint to a column:
ALTER TABLE distributors ALTER COLUMN street SET NOT NULL;

To remove a not-null constraint from a column:
ALTER TABLE distributors ALTER COLUMN street DROP NOT NULL;

To add a check constraint to a table and all its children:
ALTER TABLE distributors ADD CONSTRAINT zipchk CHECK (char_length(zipcode) = 5);

To add a check constraint only to a table and not to its children:
ALTER TABLE distributors ADD CONSTRAINT zipchk CHECK (char_length(zipcode) = 5) NO INHERIT;
(The check constraint will not be inherited by future children, either.)

To remove a check constraint from a table and all its children:
ALTER TABLE distributors DROP CONSTRAINT zipchk;

To remove a check constraint from one table only:
ALTER TABLE ONLY distributors DROP CONSTRAINT zipchk;
ALTER TABLE

(The check constraint remains in place for any child tables.)
To add a foreign key constraint to a table:

```
ALTER TABLE distributors ADD CONSTRAINT distfk FOREIGN KEY (address) REFERENCES addresses (address);
```

To add a foreign key constraint to a table with the least impact on other work:

```
ALTER TABLE distributors ADD CONSTRAINT distfk FOREIGN KEY (address) REFERENCES addresses (address) NOT VALID;
ALTER TABLE distributors VALIDATE CONSTRAINT distfk;
```

To add a (multicolumn) unique constraint to a table:

```
ALTER TABLE distributors ADD CONSTRAINT dist_id_zipcode_key UNIQUE (dist_id, zipcode);
```

To add an automatically named primary key constraint to a table, noting that a table can only ever have one primary key:

```
ALTER TABLE distributors ADD PRIMARY KEY (dist_id);
```

To move a table to a different tablespace:

```
ALTER TABLE distributors SET TABLESPACE fasttablespace;
```

To move a table to a different schema:

```
ALTER TABLE myschema.distributors SET SCHEMA yourschema;
```

To recreate a primary key constraint, without blocking updates while the index is rebuilt:

```
CREATE UNIQUE INDEX CONCURRENTLY dist_id_temp_idx ON distributors (dist_id);
ALTER TABLE distributors DROP CONSTRAINT distributors_pkey,
   ADD CONSTRAINT distributors_pkey PRIMARY KEY USING INDEX dist_id_temp_idx;
```

Compatibility

The forms `ADD` (without `USING INDEX`), `DROP`, `SET DEFAULT`, and `SET DATA TYPE` (without `USING`) conform with the SQL standard. The other forms are PostgreSQL extensions of the SQL standard. Also, the ability to specify more than one manipulation in a single `ALTER TABLE` command is an extension.

`ALTER TABLE DROP COLUMN` can be used to drop the only column of a table, leaving a zero-column table. This is an extension of SQL, which disallows zero-column tables.
See Also

CREATE TABLE
**ALTER TABLESPACE**

**Name**

ALTER TABLESPACE — change the definition of a tablespace

**Synopsis**

ALTER TABLESPACE name RENAME TO new_name
ALTER TABLESPACE name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER TABLESPACE name SET ( tablespace_option = value [, ... ] )
ALTER TABLESPACE name RESET ( tablespace_option [, ... ] )

**Description**

ALTER TABLESPACE can be used to change the definition of a tablespace.

You must own the tablespace to change the definition of a tablespace. To alter the owner, you must also be a direct or indirect member of the new owning role. (Note that superusers have these privileges automatically.)

**Parameters**

*name*

The name of an existing tablespace.

*new_name*

The new name of the tablespace. The new name cannot begin with `pg_`, as such names are reserved for system tablespaces.

*new_owner*

The new owner of the tablespace.

*tablespace_option*

A tablespace parameter to be set or reset. Currently, the only available parameters are `seq_page_cost`, `random_page_cost` and `effective_io_concurrency`. Setting either value for a particular tablespace will override the planner's usual estimate of the cost of reading pages from tables in that tablespace, as established by the configuration parameters of the same name (see `seq_page_cost`, `random_page_cost`, `effective_io_concurrency`). This may be useful if one tablespace is located on a disk which is faster or slower than the remainder of the I/O subsystem.

**Examples**

Rename tablespace `index_space` to `fast_raid`:

```
ALTER TABLESPACE index_space RENAME TO fast_raid;
```
Change the owner of tablespace `index_space`:

```
ALTER TABLESPACE index_space OWNER TO mary;
```

**Compatibility**

There is no `ALTER TABLESPACE` statement in the SQL standard.

**See Also**

`CREATE TABLESPACE`, `DROP TABLESPACE`
ALTER TEXT SEARCH CONFIGURATION

Name
ALTER TEXT SEARCH CONFIGURATION — change the definition of a text search configuration

Synopsis
ALTER TEXT SEARCH CONFIGURATION name
ADD MAPPING FOR token_type [, ... ] WITH dictionary_name [, ... ]
ALTER TEXT SEARCH CONFIGURATION name
ALTER MAPPING FOR token_type [, ... ] WITH dictionary_name [, ... ]
ALTER TEXT SEARCH CONFIGURATION name
ALTER MAPPING REPLACE old_dictionary WITH new_dictionary
ALTER TEXT SEARCH CONFIGURATION name
ALTER MAPPING FOR token_type [, ... ] REPLACE old_dictionary WITH new_dictionary
ALTER TEXT SEARCH CONFIGURATION name
DROP MAPPING [ IF EXISTS ] FOR token_type [, ... ]
ALTER TEXT SEARCH CONFIGURATION name RENAME TO new_name
ALTER TEXT SEARCH CONFIGURATION name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER TEXT SEARCH CONFIGURATION name SET SCHEMA new_schema

Description
ALTER TEXT SEARCH CONFIGURATION changes the definition of a text search configuration. You can modify its mappings from token types to dictionaries, or change the configuration’s name or owner.

You must be the owner of the configuration to use ALTER TEXT SEARCH CONFIGURATION.

Parameters

name
The name (optionally schema-qualified) of an existing text search configuration.

token_type
The name of a token type that is emitted by the configuration’s parser.

dictionary_name
The name of a text search dictionary to be consulted for the specified token type(s). If multiple dictionaries are listed, they are consulted in the specified order.

old_dictionary
The name of a text search dictionary to be replaced in the mapping.

new_dictionary
The name of a text search dictionary to be substituted for old_dictionary.

new_name
The new name of the text search configuration.
new_owner

The new owner of the text search configuration.

new_schema

The new schema for the text search configuration.

The ADD MAPPING FOR form installs a list of dictionaries to be consulted for the specified token type(s); it is an error if there is already a mapping for any of the token types. The ALTER MAPPING FOR form does the same, but first removing any existing mapping for those token types. The ALTER MAPPING REPLACE forms substitute new_dictionary for old_dictionary anywhere the latter appears. This is done for only the specified token types when FOR appears, or for all mappings of the configuration when it doesn’t. The DROP MAPPING form removes all dictionaries for the specified token type(s), causing tokens of those types to be ignored by the text search configuration. It is an error if there is no mapping for the token types, unless IF EXISTS appears.

Examples

The following example replaces the english dictionary with the swedish dictionary anywhere that english is used within my_config.

ALTER TEXT SEARCH CONFIGURATION my_config
   ALTER MAPPING REPLACE english WITH swedish;

Compatibility

There is no ALTER TEXT SEARCH CONFIGURATION statement in the SQL standard.

See Also

CREATE TEXT SEARCH CONFIGURATION, DROP TEXT SEARCH CONFIGURATION
ALTER TEXT SEARCH DICTIONARY

Name

ALTER TEXT SEARCH DICTIONARY — change the definition of a text search dictionary

Synopsis

ALTER TEXT SEARCH DICTIONARY name (  
    option [ = value ] [, ... ]  
)  
ALTER TEXT SEARCH DICTIONARY name RENAME TO new_name  
ALTER TEXT SEARCH DICTIONARY name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }  
ALTER TEXT SEARCH DICTIONARY name SET SCHEMA new_schema

Description

ALTER TEXT SEARCH DICTIONARY changes the definition of a text search dictionary. You can change the dictionary’s template-specific options, or change the dictionary’s name or owner.

You must be the owner of the dictionary to use ALTER TEXT SEARCH DICTIONARY.

Parameters

name

The name (optionally schema-qualified) of an existing text search dictionary.

option

The name of a template-specific option to be set for this dictionary.

value

The new value to use for a template-specific option. If the equal sign and value are omitted, then any previous setting for the option is removed from the dictionary, allowing the default to be used.

new_name

The new name of the text search dictionary.

new_owner

The new owner of the text search dictionary.

new_schema

The new schema for the text search dictionary.

Template-specific options can appear in any order.
Examples

The following example command changes the stopword list for a Snowball-based dictionary. Other parameters remain unchanged.

ALTER TEXT SEARCH DICTIONARY my_dict ( StopWords = newrussian );

The following example command changes the language option to dutch, and removes the stopword option entirely.

ALTER TEXT SEARCH DICTIONARY my_dict ( language = dutch, StopWords );

The following example command “updates” the dictionary’s definition without actually changing anything.

ALTER TEXT SEARCH DICTIONARY my_dict ( dummy );

(The reason this works is that the option removal code doesn’t complain if there is no such option.) This trick is useful when changing configuration files for the dictionary: the ALTER will force existing database sessions to re-read the configuration files, which otherwise they would never do if they had read them earlier.

Compatibility

There is no ALTER TEXT SEARCH DICTIONARY statement in the SQL standard.

See Also

CREATE TEXT SEARCH DICTIONARY, DROP TEXT SEARCH DICTIONARY
ALTER TEXT SEARCH PARSER

Name
ALTER TEXT SEARCH PARSER — change the definition of a text search parser

Synopsis
ALTER TEXT SEARCH PARSER name RENAME TO new_name
ALTER TEXT SEARCH PARSER name SET SCHEMA new_schema

Description
ALTER TEXT SEARCH PARSER changes the definition of a text search parser. Currently, the only supported functionality is to change the parser’s name.

You must be a superuser to use ALTER TEXT SEARCH PARSER.

Parameters

name
The name (optionally schema-qualified) of an existing text search parser.

new_name
The new name of the text search parser.

new_schema
The new schema for the text search parser.

Compatibility
There is no ALTER TEXT SEARCH PARSER statement in the SQL standard.

See Also
CREATE TEXT SEARCH PARSER, DROP TEXT SEARCH PARSER
ALTER TEXT SEARCH TEMPLATE

Name
ALTER TEXT SEARCH TEMPLATE — change the definition of a text search template

Synopsis
ALTER TEXT SEARCH TEMPLATE name RENAME TO new_name
ALTER TEXT SEARCH TEMPLATE name SET SCHEMA new_schema

Description
ALTER TEXT SEARCH TEMPLATE changes the definition of a text search template. Currently, the only supported functionality is to change the template’s name.

You must be a superuser to use ALTER TEXT SEARCH TEMPLATE.

Parameters

name
   The name (optionally schema-qualified) of an existing text search template.

new_name
   The new name of the text search template.

new_schema
   The new schema for the text search template.

Compatibility

There is no ALTER TEXT SEARCH TEMPLATE statement in the SQL standard.

See Also

CREATE TEXT SEARCH TEMPLATE, DROP TEXT SEARCH TEMPLATE
ALTER TRIGGER

**Name**

ALTER TRIGGER — change the definition of a trigger

**Synopsis**

ALTER TRIGGER name ON table_name RENAME TO new_name
ALTER TRIGGER name ON table_name DEPENDS ON EXTENSION extension_name

**Description**

ALTER TRIGGER changes properties of an existing trigger. The RENAME clause changes the name of the given trigger without otherwise changing the trigger definition. The DEPENDS ON EXTENSION clause marks the trigger as dependent on an extension, such that if the extension is dropped, the trigger will automatically be dropped as well.

You must own the table on which the trigger acts to be allowed to change its properties.

**Parameters**

*name*

The name of an existing trigger to alter.

*table_name*

The name of the table on which this trigger acts.

*new_name*

The new name for the trigger.

*extension_name*

The name of the extension that the trigger is to depend on.

**Notes**

The ability to temporarily enable or disable a trigger is provided by ALTER TABLE, not by ALTER TRIGGER, because ALTER TRIGGER has no convenient way to express the option of enabling or disabling all of a table’s triggers at once.

**Examples**

To rename an existing trigger:

ALTER TRIGGER emp_stamp ON emp RENAME TO emp_track_chgs;
To mark a trigger as being dependent on an extension:

```
ALTER TRIGGER emp_stamp ON emp DEPENDS ON EXTENSION emplib;
```

Compatibility

**ALTER TRIGGER** is a PostgreSQL extension of the SQL standard.

See Also

**ALTER TABLE**
ALTER TYPE

Name

ALTER TYPE — change the definition of a type

Synopsis

ALTER TYPE name action [, ... ]
ALTER TYPE name OWNER TO ( new_owner | CURRENT_USER | SESSION_USER )
ALTER TYPE name RENAME ATTRIBUTE attribute_name TO new_attribute_name [ CASCADE | RESTRICT ]
ALTER TYPE name RENAME TO new_name
ALTER TYPE name SET SCHEMA new_schema
ALTER TYPE name ADD VALUE [ IF NOT EXISTS ] new_enum_value [ { BEFORE | AFTER } existing_enum_value ]

where action is one of:

ADD ATTRIBUTE attribute_name data_type [ COLLATE collation ] [ CASCADE | RESTRICT ]
DROP ATTRIBUTE [ IF EXISTS ] attribute_name [ CASCADE | RESTRICT ]
ALTER ATTRIBUTE attribute_name [ SET DATA ] TYPE data_type [ COLLATE collation ] [ CASCADE | RESTRICT ]

Description

ALTER TYPE changes the definition of an existing type. There are several subforms:

ADD ATTRIBUTE

This form adds a new attribute to a composite type, using the same syntax as CREATE TYPE.

DROP ATTRIBUTE [ IF EXISTS ]

This form drops an attribute from a composite type. If IF EXISTS is specified and the attribute does not exist, no error is thrown. In this case a notice is issued instead.

SET DATA TYPE

This form changes the type of an attribute of a composite type.

OWNER

This form changes the owner of the type.

RENAME

This form changes the name of the type or the name of an individual attribute of a composite type.

SET SCHEMA

This form moves the type into another schema.

ADD VALUE [ IF NOT EXISTS ] [ BEORE | AFTER ]

This form adds a new value to an enum type. The new value’s place in the enum’s ordering can be specified as being BEFORE or AFTER one of the existing values. Otherwise, the new item is added at the end of the list of values.
ALTER TYPE

If `IF NOT EXISTS` is specified, it is not an error if the type already contains the new value: a notice is issued but no other action is taken. Otherwise, an error will occur if the new value is already present.

CASCADE

Automatically propagate the operation to typed tables of the type being altered, and their descendants.

RESTRICT

Refuse the operation if the type being altered is the type of a typed table. This is the default.

The `ADD ATTRIBUTE`, `DROP ATTRIBUTE`, and `ALTER ATTRIBUTE` actions can be combined into a list of multiple alterations to apply in parallel. For example, it is possible to add several attributes and/or alter the type of several attributes in a single command.

You must own the type to use `ALTER TYPE`. To change the schema of a type, you must also have `CREATE` privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have `CREATE` privilege on the type’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the type. However, a superuser can alter ownership of any type anyway.) To add an attribute or alter an attribute type, you must also have `USAGE` privilege on the data type.

Parameters

- **name**
  The name (possibly schema-qualified) of an existing type to alter.

- **new_name**
  The new name for the type.

- **new_owner**
  The user name of the new owner of the type.

- **new_schema**
  The new schema for the type.

- **attribute_name**
  The name of the attribute to add, alter, or drop.

- **new_attribute_name**
  The new name of the attribute to be renamed.

- **data_type**
  The data type of the attribute to add, or the new type of the attribute to alter.

- **new_enum_value**
  The new value to be added to an enum type’s list of values. Like all enum literals, it needs to be quoted.
existing_enum_value

The existing enum value that the new value should be added immediately before or after in the
denum type’s sort ordering. Like all enum literals, it needs to be quoted.

Notes

ALTER TYPE ... ADD VALUE (the form that adds a new value to an enum type) cannot be executed
inside a transaction block.

Comparisons involving an added enum value will sometimes be slower than comparisons involving
only original members of the enum type. This will usually only occur if BEFORE or AFTER is used to
set the new value’s sort position somewhere other than at the end of the list. However, sometimes it
will happen even though the new value is added at the end (this occurs if the OID counter “wrapped
around” since the original creation of the enum type). The slowdown is usually insignificant; but if
it matters, optimal performance can be regained by dropping and recreating the enum type, or by
dumping and reloading the database.

Examples

To rename a data type:

ALTER TYPE electronic_mail RENAME TO email;

To change the owner of the type email to joe:

ALTER TYPE email OWNER TO joe;

To change the schema of the type email to customers:

ALTER TYPE email SET SCHEMA customers;

To add a new attribute to a type:

ALTER TYPE compfoo ADD ATTRIBUTE f3 int;

To add a new value to an enum type in a particular sort position:

ALTER TYPE colors ADD VALUE 'orange' AFTER 'red';

Compatibility

The variants to add and drop attributes are part of the SQL standard; the other variants are PostgreSQL
extensions.
See Also

CREATE TYPE, DROP TYPE
ALTER USER

Name
ALTER USER — change a database role

Synopsis
ALTER USER role_specification [ WITH ] option [ ... ]

where option can be:

- SUPERUSER | NOSUPERUSER
- CREATEDB | NOCREATEDB
- CREATETABLE | NOCREATETABLE
- INHERIT | NOINHERIT
- LOGIN | NOLOGIN
- REPLICATION | NOREPLICATION
- BYPASSRLS | NOBYPASSRLS
- CONNECTION LIMIT connlimit
- [ ENCRYPTED | UNENCRYPTED ] PASSWORD 'password'
- VALID UNTIL 'timestamp'

ALTER USER name RENAME TO new_name

ALTER USER { role_specification | ALL } [ IN DATABASE database_name ] SET configuration_parameter

ALTER USER { role_specification | ALL } [ IN DATABASE database_name ] SET configuration_parameter FROM CURRENT

ALTER USER { role_specification | ALL } [ IN DATABASE database_name ] RESET configuration_parameter

ALTER USER { role_specification | ALL } [ IN DATABASE database_name ] RESET ALL

where role_specification can be:

- role_name
- CURRENT_USER
- SESSION_USER

Description
ALTER USER is now an alias for ALTER ROLE.

Compatibility
The ALTER USER statement is a PostgreSQL extension. The SQL standard leaves the definition of users to the implementation.

See Also
ALTER ROLE
ALTER USER MAPPING

Name
ALTER USER MAPPING — change the definition of a user mapping

Synopsis
ALTER USER MAPPING FOR { user_name | USER | CURRENT_USER | SESSION_USER | PUBLIC }
    SERVER server_name
    OPTIONS ( [ ADD | SET | DROP ] option [‘value’] [, ... ] )

Description
ALTER USER MAPPING changes the definition of a user mapping.
The owner of a foreign server can alter user mappings for that server for any user. Also, a user can alter a user mapping for their own user name if USAGE privilege on the server has been granted to the user.

Parameters
user_name
    User name of the mapping. CURRENT_USER and USER match the name of the current user. PUBLIC is used to match all present and future user names in the system.

server_name
    Server name of the user mapping.
    OPTIONS ( [ ADD | SET | DROP ] option [‘value’] [, ... ] )
    Change options for the user mapping. The new options override any previously specified options. ADD, SET, and DROP specify the action to be performed. ADD is assumed if no operation is explicitly specified. Option names must be unique; options are also validated by the server’s foreign-data wrapper.

Examples
Change the password for user mapping bob, server foo:

ALTER USER MAPPING FOR bob SERVER foo OPTIONS (SET password ‘public’);

Compatibility
ALTER USER MAPPING conforms to ISO/IEC 9075-9 (SQL/MED). There is a subtle syntax issue: The standard omits the FOR key word. Since both CREATE USER MAPPING and DROP USER
ALTER USER MAPPING

MAPPING use FOR in analogous positions, and IBM DB2 (being the other major SQL/MED implementation) also requires it for ALTER USER MAPPING, PostgreSQL diverges from the standard here in the interest of consistency and interoperability.

See Also

CREATE USER MAPPING, DROP USER MAPPING
**ALTER VIEW**

**Name**
ALTER VIEW — change the definition of a view

**Synopsis**

ALTER VIEW [ IF EXISTS ] name ALTER [ COLUMN ] column_name SET DEFAULT expression
ALTER VIEW [ IF EXISTS ] name ALTER [ COLUMN ] column_name DROP DEFAULT
ALTER VIEW [ IF EXISTS ] name OWNER TO { new_owner | CURRENT_USER | SESSION_USER }
ALTER VIEW [ IF EXISTS ] name RENAME TO new_name
ALTER VIEW [ IF EXISTS ] name SET SCHEMA new_schema
ALTER VIEW [ IF EXISTS ] name SET ( view_option_name [= view_option_value] [, ... ] )
ALTER VIEW [ IF EXISTS ] name RESET ( view_option_name [, ... ] )

**Description**

ALTER VIEW changes various auxiliary properties of a view. (If you want to modify the view’s defining query, use CREATE OR REPLACE VIEW.)

You must own the view to use ALTER VIEW. To change a view’s schema, you must also have CREATE privilege on the new schema. To alter the owner, you must also be a direct or indirect member of the new owning role, and that role must have CREATE privilege on the view’s schema. (These restrictions enforce that altering the owner doesn’t do anything you couldn’t do by dropping and recreating the view. However, a superuser can alter ownership of any view anyway.)

**Parameters**

*name*

The name (optionally schema-qualified) of an existing view.

*IF EXISTS*

Do not throw an error if the view does not exist. A notice is issued in this case.

*SET/DROP DEFAULT*

These forms set or remove the default value for a column. A view column’s default value is substituted into any INSERT or UPDATE command whose target is the view, before applying any rules or triggers for the view. The view’s default will therefore take precedence over any default values from underlying relations.

*new_owner*

The user name of the new owner of the view.

*new_name*

The new name for the view.

*new_schema*

The new schema for the view.
SET ( view_option_name [= view_option_value] [, ... ] )
RESET ( view_option_name [, ... ] )

Sets or resets a view option. Currently supported options are:

check_option (string)
  Changes the check option of the view. The value must be local or cascaded.

security_barrier (boolean)
  Changes the security-barrier property of the view. The value must be Boolean value, such as true or false.

Notes
For historical reasons, ALTER TABLE can be used with views too; but the only variants of ALTER TABLE that are allowed with views are equivalent to the ones shown above.

Examples
To rename the view foo to bar:

ALTER VIEW foo RENAME TO bar;

To attach a default column value to an updatable view:

CREATE TABLE base_table (id int, ts timestamptz);
CREATE VIEW a_view AS SELECT * FROM base_table;
ALTER VIEW a_view ALTER COLUMN ts SET DEFAULT now();
INSERT INTO base_table(id) VALUES(1); -- ts will receive a NULL
INSERT INTO a_view(id) VALUES(2); -- ts will receive the current time

Compatibility
ALTER VIEW is a PostgreSQL extension of the SQL standard.

See Also
CREATE VIEW, DROP VIEW
ANALYZE

Name

ANALYZE — collect statistics about a database

Synopsis

ANALYZE [ VERBOSE ] [ table_name [ ( column_name [, ...] ) ] ]

Description

ANALYZE collects statistics about the contents of tables in the database, and stores the results in the \texttt{pg\_statistic} system catalog. Subsequently, the query planner uses these statistics to help determine the most efficient execution plans for queries.

With no parameter, \texttt{ANALYZE} examines every table in the current database. With a parameter, \texttt{ANALYZE} examines only that table. It is further possible to give a list of column names, in which case only the statistics for those columns are collected.

Parameters

\texttt{VERBOSE}

Enables display of progress messages.

\texttt{table\_name}

The name (possibly schema-qualified) of a specific table to analyze. If omitted, all regular tables (but not foreign tables) in the current database are analyzed.

\texttt{column\_name}

The name of a specific column to analyze. Defaults to all columns.

Outputs

When \texttt{VERBOSE} is specified, \texttt{ANALYZE} emits progress messages to indicate which table is currently being processed. Various statistics about the tables are printed as well.

Notes

To analyze a table, one must ordinarily be the table’s owner or a superuser. However, database owners are allowed to analyze all tables in their databases, except shared catalogs. (The restriction for shared catalogs means that a true database-wide \texttt{ANALYZE} can only be performed by a superuser.) \texttt{ANALYZE} will skip over any tables that the calling user does not have permission to analyze.

Foreign tables are analyzed only when explicitly selected. Not all foreign data wrappers support \texttt{ANALYZE}. If the table’s wrapper does not support \texttt{ANALYZE}, the command prints a warning and does nothing.
In the default PostgreSQL configuration, the autovacuum daemon (see Section 24.1.6) takes care of automatic analyzing of tables when they are first loaded with data, and as they change throughout regular operation. When autovacuum is disabled, it is a good idea to run `ANALYZE` periodically, or just after making major changes in the contents of a table. Accurate statistics will help the planner to choose the most appropriate query plan, and thereby improve the speed of query processing. A common strategy for read-mostly databases is to run `VACUUM` and `ANALYZE` once a day during a low-usage time of day. (This will not be sufficient if there is heavy update activity.)

`ANALYZE` requires only a read lock on the target table, so it can run in parallel with other activity on the table.

The statistics collected by `ANALYZE` usually include a list of some of the most common values in each column and a histogram showing the approximate data distribution in each column. One or both of these can be omitted if `ANALYZE` deems them uninteresting (for example, in a unique-key column, there are no common values) or if the column data type does not support the appropriate operators. There is more information about the statistics in Chapter 24.

For large tables, `ANALYZE` takes a random sample of the table contents, rather than examining every row. This allows even very large tables to be analyzed in a small amount of time. Note, however, that the statistics are only approximate, and will change slightly each time `ANALYZE` is run, even if the actual table contents did not change. This might result in small changes in the planner’s estimated costs shown by `EXPLAIN`. In rare situations, this non-determinism will cause the planner’s choices of query plans to change after `ANALYZE` is run. To avoid this, raise the amount of statistics collected by `ANALYZE`, as described below.

The extent of analysis can be controlled by adjusting the `default_statistics_target` configuration variable, or on a column-by-column basis by setting the per-column statistics target with `ALTER TABLE ... ALTER COLUMN ... SET STATISTICS` (see `ALTER TABLE`). The target value sets the maximum number of entries in the most-common-value list and the maximum number of bins in the histogram. The default target value is 100, but this can be adjusted up or down to trade off accuracy of planner estimates against the time taken for `ANALYZE` and the amount of space occupied in `pg_statistic`. In particular, setting the statistics target to zero disables collection of statistics for that column. It might be useful to do that for columns that are never used as part of the `WHERE`, `GROUP BY`, or `ORDER BY` clauses of queries, since the planner will have no use for statistics on such columns.

The largest statistics target among the columns being analyzed determines the number of table rows sampled to prepare the statistics. Increasing the target causes a proportional increase in the time and space needed to do `ANALYZE`.

One of the values estimated by `ANALYZE` is the number of distinct values that appear in each column. Because only a subset of the rows are examined, this estimate can sometimes be quite inaccurate, even with the largest possible statistics target. If this inaccuracy leads to bad query plans, a more accurate value can be determined manually and then installed with `ALTER TABLE ... ALTER COLUMN ... SET (n_distinct = ...)` (see `ALTER TABLE`).

If the table being analyzed has one or more children, `ANALYZE` will gather statistics twice: once on the rows of the parent table only, and a second time on the rows of the parent table with all of its children. This second set of statistics is needed when planning queries that traverse the entire inheritance tree. The autovacuum daemon, however, will only consider inserts or updates on the parent table itself when deciding whether to trigger an automatic analyze for that table. If that table is rarely inserted into or updated, the inheritance statistics will not be up to date unless you run `ANALYZE` manually.

If any of the child tables are foreign tables whose foreign data wrappers do not support `ANALYZE`, those child tables are ignored while gathering inheritance statistics.

If the table being analyzed is completely empty, `ANALYZE` will not record new statistics for that table. Any existing statistics will be retained.
Compatibility

There is no `ANALYZE` statement in the SQL standard.

See Also

VACUUM, vacuumdb, Section 19.4.4, Section 24.1.6
BEGIN

Name
BEGIN — start a transaction block

Synopsis
BEGIN [ WORK | TRANSACTION ] [ transaction_mode [, ...] ]

where transaction_mode is one of:

    ISOLATION LEVEL | SERIALIZABLE | REPEATABLE READ | READ COMMITTED | READ UNCOMMITTED
    READ WRITE | READ ONLY
    [ NOT ] DEFERRABLE

Description
BEGIN initiates a transaction block, that is, all statements after a BEGIN command will be executed in a single transaction until an explicit COMMIT or ROLLBACK is given. By default (without BEGIN), PostgreSQL executes transactions in “autocommit” mode, that is, each statement is executed in its own transaction and a commit is implicitly performed at the end of the statement (if execution was successful, otherwise a rollback is done).

Statements are executed more quickly in a transaction block, because transaction start/commit requires significant CPU and disk activity. Execution of multiple statements inside a transaction is also useful to ensure consistency when making several related changes: other sessions will be unable to see the intermediate states wherein not all the related updates have been done.

If the isolation level, read/write mode, or deferrable mode is specified, the new transaction has those characteristics, as if SET TRANSACTION was executed.

Parameters
WORK
TRANSACTION

Optional key words. They have no effect.

Refer to SET TRANSACTION for information on the meaning of the other parameters to this statement.

Notes
START TRANSACTION has the same functionality as BEGIN.

Use COMMIT or ROLLBACK to terminate a transaction block.

Issuing BEGIN when already inside a transaction block will provoke a warning message. The state of the transaction is not affected. To nest transactions within a transaction block, use savepoints (see SAVEPOINT).
For reasons of backwards compatibility, the commas between successive `transaction_modes` can be omitted.

**Examples**

To begin a transaction block:

```
BEGIN;
```

**Compatibility**

`BEGIN` is a PostgreSQL language extension. It is equivalent to the SQL-standard command `START TRANSACTION`, whose reference page contains additional compatibility information.

The `DEFERRABLE transaction_mode` is a PostgreSQL language extension.

Incidentally, the `BEGIN` key word is used for a different purpose in embedded SQL. You are advised to be careful about the transaction semantics when porting database applications.

**See Also**

`COMMIT`, `ROLLBACK`, `START TRANSACTION`, `SAVEPOINT`
CHECKPOINT

Name

CHECKPOINT — force a transaction log checkpoint

Synopsis

CHECKPOINT

Description

A checkpoint is a point in the transaction log sequence at which all data files have been updated to reflect the information in the log. All data files will be flushed to disk. Refer to Section 30.4 for more details about what happens during a checkpoint.

The CHECKPOINT command forces an immediate checkpoint when the command is issued, without waiting for a regular checkpoint scheduled by the system (controlled by the settings in Section 19.5.2). CHECKPOINT is not intended for use during normal operation.

If executed during recovery, the CHECKPOINT command will force a restartpoint (see Section 30.4) rather than writing a new checkpoint.

Only superusers can call CHECKPOINT.

Compatibility

The CHECKPOINT command is a PostgreSQL language extension.
CLOSE

Name
CLOSE — close a cursor

Synopsis
CLOSE { name | ALL }

Description
CLOSE frees the resources associated with an open cursor. After the cursor is closed, no subsequent
operations are allowed on it. A cursor should be closed when it is no longer needed.

Every non-holdable open cursor is implicitly closed when a transaction is terminated by COMMIT
or ROLLBACK. A holdable cursor is implicitly closed if the transaction that created it aborts via
ROLLBACK. If the creating transaction successfully commits, the holdable cursor remains open un-
til an explicit CLOSE is executed, or the client disconnects.

Parameters

name
The name of an open cursor to close.

ALL
Close all open cursors.

Notes
PostgreSQL does not have an explicit OPEN cursor statement; a cursor is considered open when it is
declared. Use the DECLARE statement to declare a cursor.

You can see all available cursors by querying the pg_cursors system view.

If a cursor is closed after a savepoint which is later rolled back, the CLOSE is not rolled back; that is,
the cursor remains closed.

Examples
Close the cursor liahona:

CLOSE liahona;
Compatibility

CLOSE is fully conforming with the SQL standard. CLOSE ALL is a PostgreSQL extension.

See Also

DECLARE, FETCH, MOVE
**CLUSTER**

**Name**

CLUSTER — cluster a table according to an index

**Synopsis**

CLUSTER [VERBOSE] table_name [ USING index_name ]
CLUSTER [VERBOSE]

**Description**

CLUSTER instructs PostgreSQL to cluster the table specified by `table_name` based on the index specified by `index_name`. The index must already have been defined on `table_name`.

When a table is clustered, it is physically reordered based on the index information. Clustering is a one-time operation: when the table is subsequently updated, the changes are not clustered. That is, no attempt is made to store new or updated rows according to their index order. (If one wishes, one can periodically recluster by issuing the command again. Also, setting the table’s `fillfactor` storage parameter to less than 100% can aid in preserving cluster ordering during updates, since updated rows are kept on the same page if enough space is available there.)

When a table is clustered, PostgreSQL remembers which index it was clustered by. The form `CLUSTER table_name` reclusters the table using the same index as before. You can also use the `CLUSTER` or `SET WITHOUT CLUSTER` forms of `ALTER TABLE` to set the index to be used for future cluster operations, or to clear any previous setting.

`CLUSTER` without any parameter reclusters all the previously-clustered tables in the current database that the calling user owns, or all such tables if called by a superuser. This form of `CLUSTER` cannot be executed inside a transaction block.

When a table is being clustered, an `ACCESS EXCLUSIVE` lock is acquired on it. This prevents any other database operations (both reads and writes) from operating on the table until the `CLUSTER` is finished.

**Parameters**

`table_name`

The name (possibly schema-qualified) of a table.

`index_name`

The name of an index.

`VERBOSE`

Prints a progress report as each table is clustered.
Notes

In cases where you are accessing single rows randomly within a table, the actual order of the data in the table is unimportant. However, if you tend to access some data more than others, and there is an index that groups them together, you will benefit from using CLUSTER. If you are requesting a range of indexed values from a table, or a single indexed value that has multiple rows that match, CLUSTER will help because once the index identifies the table page for the first row that matches, all other rows that match are probably already on the same table page, and so you save disk accesses and speed up the query.

CLUSTER can re-sort the table using either an index scan on the specified index, or (if the index is a b-tree) a sequential scan followed by sorting. It will attempt to choose the method that will be faster, based on planner cost parameters and available statistical information.

When an index scan is used, a temporary copy of the table is created that contains the table data in the index order. Temporary copies of each index on the table are created as well. Therefore, you need free space on disk at least equal to the sum of the table size and the index sizes.

When a sequential scan and sort is used, a temporary sort file is also created, so that the peak temporary space requirement is as much as double the table size, plus the index sizes. This method is often faster than the index scan method, but if the disk space requirement is intolerable, you can disable this choice by temporarily setting enable_sort to off.

It is advisable to set maintenance_work_mem to a reasonably large value (but not more than the amount of RAM you can dedicate to the CLUSTER operation) before clustering.

Because the planner records statistics about the ordering of tables, it is advisable to run ANALYZE on the newly clustered table. Otherwise, the planner might make poor choices of query plans.

Because CLUSTER remembers which indexes are clustered, one can cluster the tables one wants clustered manually the first time, then set up a periodic maintenance script that executes CLUSTER without any parameters, so that the desired tables are periodically reclustered.

Examples

Cluster the table employees on the basis of its index employees_ind:

CLUSTER employees USING employees_ind;

Cluster the employees table using the same index that was used before:

CLUSTER employees;

Cluster all tables in the database that have previously been clustered:

CLUSTER;

Compatibility

There is no CLUSTER statement in the SQL standard.
The syntax

`CLUSTER index_name ON table_name`

is also supported for compatibility with pre-8.3 PostgreSQL versions.

See Also

`clusterdb`
**Name**

COMMENT — define or change the comment of an object

**Synopsis**

COMMENT ON

{ ACCESS METHOD object_name |
  AGGREGATE aggregate_name ( aggregate_signature ) |
  CAST (source_type AS target_type) |
  COLLATION object_name |
  COLUMN relation_name.column_name |
  CONSTRAINT constraint_name ON table_name |
  CONSTRAINT constraint_name ON DOMAIN domain_name |
  CONVERSION object_name |
  DATABASE object_name |
  DOMAIN object_name |
  EXTENSION object_name |
  EVENT TRIGGER object_name |
  FOREIGN DATA WRAPPER object_name |
  FOREIGN TABLE object_name |
  FUNCTION function_name ( [ [ argmode ] [ argname ] argtype [, ... ] ] ) |
  INDEX object_name |
  LARGE OBJECT large_object_oid |
  MATERIALIZED VIEW object_name |
  OPERATOR operator_name (left_type, right_type) |
  OPERATOR CLASS object_name USING index_method |
  OPERATOR FAMILY object_name USING index_method |
  POLICY policy_name ON table_name |
  [ PROCEDURAL ] LANGUAGE object_name |
  ROLE object_name |
  RULE rule_name ON table_name |
  SCHEMA object_name |
  SEQUENCE object_name |
  SERVER object_name |
  TABLE object_name |
  TABLESPACE object_name |
  TEXT SEARCH CONFIGURATION object_name |
  TEXT SEARCH DICTIONARY object_name |
  TEXT SEARCH PARSER object_name |
  TEXT SEARCH TEMPLATE object_name |
  TRANSFORM FOR type_name LANGUAGE lang_name |
  TRIGGER trigger_name ON table_name |
  TYPE object_name |
  VIEW object_name
}

where aggregate_signature is:

* |
[ argmode ] [ argname ] argtype [, ... ] |
[ [ argmode ] [ argname ] argtype [, ... ] ] ORDER BY [ argmode ] [ argname ] argtype [, ... ] |
**Description**

`COMMENT` stores a comment about a database object. Only one comment string is stored for each object, so to modify a comment, issue a new `COMMENT` command for the same object. To remove a comment, write `NULL` in place of the text string. Comments are automatically dropped when their object is dropped.

For most kinds of object, only the object’s owner can set the comment. Roles don’t have owners, so the rule for `COMMENT ON ROLE` is that you must be superuser to comment on a superuser role, or have the `CREATEROLE` privilege to comment on non-superuser roles. Likewise, access methods don’t have owners either; you must be superuser to comment on an access method. Of course, a superuser can comment on anything.

Comments can be viewed using psql’s `\d` family of commands. Other user interfaces to retrieve comments can be built atop the same built-in functions that psql uses, namely `obj_description`, `col_description`, and `shobj_description` (see Table 9-67).

**Parameters**

- `object_name`
- `relation_name.column_name`
- `aggregate_name`
- `constraint_name`
- `function_name`
- `operator_name`
- `policy_name`
- `rule_name`
- `trigger_name`

The name of the object to be commented. Names of tables, aggregates, collations, conversions, domains, foreign tables, functions, indexes, operators, operator classes, operator families, sequences, text search objects, types, and views can be schema-qualified. When commenting on a column, `relation_name` must refer to a table, view, composite type, or foreign table.

- `table_name`
- `domain_name`

When creating a comment on a constraint, a trigger, a rule or a policy these parameters specify the name of the table or domain on which that object is defined.

- `source_type`

  The name of the source data type of the cast.

- `target_type`

  The name of the target data type of the cast.

- `argmode`

  The mode of a function or aggregate argument: `IN`, `OUT`, `INOUT`, or `VARIADIC`. If omitted, the default is `IN`. Note that `COMMENT` does not actually pay any attention to `OUT` arguments, since only the input arguments are needed to determine the function’s identity. So it is sufficient to list the `IN`, `INOUT`, and `VARIADIC` arguments.
argname

The name of a function or aggregate argument. Note that `COMMENT` does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

argtype

The data type of a function or aggregate argument.

large_object_oid

The OID of the large object.

left_type
	right_type

The data type(s) of the operator’s arguments (optionally schema-qualified). Write `NONE` for the missing argument of a prefix or postfix operator.

PROCEDURAL

This is a noise word.

type_name

The name of the data type of the transform.

lang_name

The name of the language of the transform.

text

The new comment, written as a string literal; or `NULL` to drop the comment.

Notes

There is presently no security mechanism for viewing comments: any user connected to a database can see all the comments for objects in that database. For shared objects such as databases, roles, and tablespaces, comments are stored globally so any user connected to any database in the cluster can see all the comments for shared objects. Therefore, don’t put security-critical information in comments.

Examples

Attach a comment to the table `mytable`:

```
COMMENT ON TABLE mytable IS 'This is my table.';
```

Remove it again:

```
COMMENT ON TABLE mytable IS NULL;
```

Some more examples:

```
COMMENT ON ACCESS METHOD rtree IS 'R-Tree access method';
COMMENT ON AGGREGATE my_aggregate (double precision) IS 'Computes sample variance';
COMMENT ON CAST (text AS int4) IS 'Allow casts from text to int4';
COMMENT ON COLLATION "fr_CA" IS 'Canadian French';
```
COMMENT ON COLUMN my_table.my_column IS 'Employee ID number';
COMMENT ON CONVERSION my_conv IS 'Conversion to UTF8';
COMMENT ON CONSTRAINT bar_col_cons ON bar IS 'Constrains column col';
COMMENT ON CONSTRAINT dom_col_constr ON DOMAIN dom IS 'Constrains col of domain';
COMMENT ON DATABASE my_database IS 'Development Database';
COMMENT ON DOMAIN my_domain IS 'Email Address Domain';
COMMENT ON EXTENSION hstore IS 'Implements the hstore data type';
COMMENT ON FOREIGN DATA WRAPPER mywrapper IS 'My foreign data wrapper';
COMMENT ON FOREIGN TABLE my_foreign_table IS 'Employee Information in other database';
COMMENT ON FUNCTION my_function (timestamp) IS 'Returns Roman Numeral';
COMMENT ON INDEX my_index IS 'Enforces uniqueness on employee ID';
COMMENT ON LANGUAGE plpython IS 'Python support for stored procedures';
COMMENT ON LARGE OBJECT 346344 IS 'Planning document';
COMMENT ON MATERIALIZED VIEW my_matview IS 'Summary of order history';
COMMENT ON OPERATOR ^ (text, text) IS 'Performs intersection of two texts';
COMMENT ON OPERATOR - (NONE, integer) IS 'Unary minus';
COMMENT ON OPERATOR CLASS int4ops USING btree IS '4 byte integer operators for btrees';
COMMENT ON OPERATOR FAMILY integer_ops USING btree IS 'All integer operators for btrees';
COMMENT ON POLICY my_policy ON mytable IS 'Filter rows by users';
COMMENT ON ROLE my_role IS 'Administration group for finance tables';
COMMENT ON RULE my_rule ON my_table IS 'Logs updates of employee records';
COMMENT ON SCHEMA my_schema IS 'Departmental data';
COMMENT ON SEQUENCE my_sequence IS 'Used to generate primary keys';
COMMENT ON SERVER myserver IS 'My foreign server';
COMMENT ON TABLE my_schema.my_table IS 'Employee Information';
COMMENT ON TABLESPACE my_tablespace IS 'Tablespace for indexes';
COMMENT ON TEXT SEARCH CONFIGURATION my_config IS 'Special word filtering';
COMMENT ON TEXT SEARCH DICTIONARY swedish IS 'Snowball stemmer for Swedish language';
COMMENT ON TEXT SEARCH PARSER my_parser IS 'Splits text into words';
COMMENT ON TEXT SEARCH TEMPLATE snowball IS 'Snowball stemmer';
COMMENT ON TRANSFORM FOR hstore LANGUAGE plpythonu IS 'Transform between hstore and Python dict';
COMMENT ON TRIGGER my_trigger ON my_table IS 'Used for RI';
COMMENT ON TYPE complex IS 'Complex number data type';
COMMENT ON VIEW my_view IS 'View of departmental costs';

Compatibility

There is no COMMENT command in the SQL standard.
COMMIT

Name
COMMIT — commit the current transaction

Synopsis
COMMIT [ WORK | TRANSACTION ]

Description
COMMIT commits the current transaction. All changes made by the transaction become visible to others and are guaranteed to be durable if a crash occurs.

Parameters
WORK
TRANSACTION
Optional key words. They have no effect.

Notes
Use ROLLBACK to abort a transaction.
Issuing COMMIT when not inside a transaction does no harm, but it will provoke a warning message.

Examples
To commit the current transaction and make all changes permanent:

COMMIT;

Compatibility
The SQL standard only specifies the two forms COMMIT and COMMIT WORK. Otherwise, this command is fully conforming.

See Also
BEGIN, ROLLBACK
COMMIT PREPARED

Name
COMMIT PREPARED — commit a transaction that was earlier prepared for two-phase commit

Synopsis
COMMIT PREPARED transaction_id

Description
COMMIT PREPARED commits a transaction that is in prepared state.

Parameters

transaction_id

The transaction identifier of the transaction that is to be committed.

Notes
To commit a prepared transaction, you must be either the same user that executed the transaction originally, or a superuser. But you do not have to be in the same session that executed the transaction.

This command cannot be executed inside a transaction block. The prepared transaction is committed immediately.

All currently available prepared transactions are listed in the pg_prepared_xacts system view.

Examples
Commit the transaction identified by the transaction identifier foobar:

COMMIT PREPARED 'foobar';

Compatibility
COMMIT PREPARED is a PostgreSQL extension. It is intended for use by external transaction management systems, some of which are covered by standards (such as X/Open XA), but the SQL side of those systems is not standardized.
See Also

PREPARE TRANSACTION, ROLLBACK PREPARED
COPY

Name
COPY — copy data between a file and a table

Synopsis
COPY table_name [ ( column_name [, ...] ) ]
FROM { ’filename’ | PROGRAM ’command’ | STDIN }
[ [ WITH ] ( option [, ...] ) ]
COPY { table_name [ ( column_name [, ...] ) ] | ( query ) } TO { ’filename’ | PROGRAM ’command’ | STDOUT }
[ [ WITH ] ( option [, ...] ) ]

where option can be one of:

- FORMAT format_name
- OIDS [ boolean ]
- FREEZE [ boolean ]
- DELIMITER ’delimiter_character’
- NULL ’null_string’
- HEADER [ boolean ]
- QUOTE ’quote_character’
- ESCAPE ’escape_character’
- FORCE_QUOTE { ( column_name [, ...] ) | * }
- FORCE_NOT_NULL ( column_name [, ...] )
- FORCE_NULL ( column_name [, ...] )
- ENCODING ’encoding_name’

Description
COPY moves data between PostgreSQL tables and standard file-system files. COPY TO copies the contents of a table to a file, while COPY FROM copies data from a file to a table (appending the data to whatever is in the table already). COPY TO can also copy the results of a SELECT query.

If a column list is specified, COPY TO copies only the data in the specified columns to the file. For COPY FROM, each field in the file is inserted, in order, into the specified column. Table columns not specified in the COPY FROM column list will receive their default values.

COPY with a file name instructs the PostgreSQL server to directly read from or write to a file. The file must be accessible by the PostgreSQL user (the user ID the server runs as) and the name must be specified from the viewpoint of the server. When PROGRAM is specified, the server executes the given command and reads from the standard output of the program, or writes to the standard input of the program. The command must be specified from the viewpoint of the server, and be executable by the PostgreSQL user. When STDIN or STDOUT is specified, data is transmitted via the connection between the client and the server.

1411
Parameters

table_name
The name (optionally schema-qualified) of an existing table.

column_name
An optional list of columns to be copied. If no column list is specified, all columns of the table will be copied.

query
A SELECT, VALUES, INSERT, UPDATE or DELETE command whose results are to be copied. Note that parentheses are required around the query. For INSERT, UPDATE and DELETE queries a RETURNING clause must be provided, and the target relation must not have a conditional rule, nor an ALSO rule, nor an INSTEAD rule that expands to multiple statements.

filename
The path name of the input or output file. An input file name can be an absolute or relative path, but an output file name must be an absolute path. Windows users might need to use an "E" string and double any backslashes used in the path name.

PROGRAM
A command to execute. In COPY FROM, the input is read from standard output of the command, and in COPY TO, the output is written to the standard input of the command. Note that the command is invoked by the shell, so if you need to pass any arguments to shell command that come from an untrusted source, you must be careful to strip or escape any special characters that might have a special meaning for the shell. For security reasons, it is best to use a fixed command string, or at least avoid passing any user input in it.

STDIN
Specifies that input comes from the client application.

STDOUT
Specifies that output goes to the client application.

boolean
Specifies whether the selected option should be turned on or off. You can write TRUE, ON, or 1 to enable the option, and FALSE, OFF, or 0 to disable it. The boolean value can also be omitted, in which case TRUE is assumed.

FORMAT
Selects the data format to be read or written: text, csv (Comma Separated Values), or binary. The default is text.

OIDS
Specifies copying the OID for each row. (An error is raised if OIDS is specified for a table that does not have OIDs, or in the case of copying a query.)

FREEZE
Requests copying the data with rows already frozen, just as they would be after running the VACUUM FREEZE command. This is intended as a performance option for initial data loading. Rows will be frozen only if the table being loaded has been created or truncated in the current
COPY

subtransaction, there are no cursors open and there are no older snapshots held by this transac-
tion.

Note that all other sessions will immediately be able to see the data once it has been successfully
loaded. This violates the normal rules of MVCC visibility and users specifying should be aware
of the potential problems this might cause.

DELIMITER

Specifies the character that separates columns within each row (line) of the file. The default is a
tab character in text format, a comma in CSV format. This must be a single one-byte character.
This option is not allowed when using binary format.

NULL

Specifies the string that represents a null value. The default is \N (backslash-N) in text format,
and an unquoted empty string in CSV format. You might prefer an empty string even in text
format for cases where you don’t want to distinguish nulls from empty strings. This option is not
allowed when using binary format.

Note: When using COPY FROM, any data item that matches this string will be stored as a null
value, so you should make sure that you use the same string as you used with COPY TO.

HEADER

Specifies that the file contains a header line with the names of each column in the file. On output,
the first line contains the column names from the table, and on input, the first line is ignored.
This option is allowed only when using CSV format.

QUOTE

Specifies the quoting character to be used when a data value is quoted. The default is double-
quote. This must be a single one-byte character. This option is allowed only when using CSV
format.

ESCAPE

Specifies the character that should appear before a data character that matches the QUOTE value.
The default is the same as the QUOTE value (so that the quoting character is doubled if it appears
in the data). This must be a single one-byte character. This option is allowed only when using
CSV format.

FORCE_QUOTE

Forces quoting to be used for all non-NULL values in each specified column. NULL output is never
quoted. If * is specified, non-NULL values will be quoted in all columns. This option is allowed
only in COPY TO, and only when using CSV format.

FORCE_NOT_NULL

Do not match the specified columns’ values against the null string. In the default case where the
null string is empty, this means that empty values will be read as zero-length strings rather than
nulls, even when they are not quoted. This option is allowed only in COPY FROM, and only when
using CSV format.

FORCE_NULL

Match the specified columns’ values against the null string, even if it has been quoted, and if
a match is found set the value to NULL. In the default case where the null string is empty, this
**COPY**

converts a quoted empty string into NULL. This option is allowed only in **COPY FROM**, and only when using **CSV** format.

**ENCODING**

Specifies that the file is encoded in the `encoding_name`. If this option is omitted, the current client encoding is used. See the Notes below for more details.

**Outputs**

On successful completion, a **COPY** command returns a command tag of the form

```
COPY count
```

The `count` is the number of rows copied.

**Note:** `psql` will print this command tag only if the command was not **COPY ... TO STDOUT**, or the equivalent `psql` meta-command `\copy ... to stdout`. This is to prevent confusing the command tag with the data that was just printed.

**Notes**

**COPY** can only be used with plain tables, not with views. However, you can write **COPY** (SELECT * FROM `viewname`) TO ... for example **COPY table TO** shows the same data as **SELECT * FROM ONLY table**. But **COPY** (SELECT * FROM `table`) TO ... can be used to dump all of the data in an inheritance hierarchy.

You must have select privilege on the table whose values are read by **COPY TO**, and insert privilege on the table into which values are inserted by **COPY FROM**. It is sufficient to have column privileges on the column(s) listed in the command.

If row-level security is enabled for the table, the relevant **SELECT** policies will apply to **COPY table TO** statements. Currently, **COPY FROM** is not supported for tables with row-level security. Use equivalent **INSERT** statements instead.

Files named in a **COPY** command are read or written directly by the server, not by the client application. Therefore, they must reside on or be accessible to the database server machine, not the client. They must be accessible to and readable or writable by the PostgreSQL user (the user ID the server runs as), not the client. Similarly, the command specified with **PROGRAM** is executed directly by the server, not by the client application, must be executable by the PostgreSQL user. **COPY** naming a file or command is only allowed to database superusers, since it allows reading or writing any file that the server has privileges to access.

Do not confuse **COPY** with the `psql` instruction `\copy`. `\copy` invokes **COPY FROM STDIN** or **COPY TO STDOUT**, and then fetches/stores the data in a file accessible to the `psql` client. Thus, file accessibility and access rights depend on the client rather than the server when `\copy` is used.

It is recommended that the file name used in **COPY** always be specified as an absolute path. This is enforced by the server in the case of **COPY TO**, but for **COPY FROM** you do have the option of reading from a file specified by a relative path. The path will be interpreted relative to the working directory of the server process (normally the cluster’s data directory), not the client’s working directory.
Executing a command with `PROGRAM` might be restricted by the operating system’s access control mechanisms, such as SELinux.

`COPY FROM` will invoke any triggers and check constraints on the destination table. However, it will not invoke rules.

`COPY` input and output is affected by `DateStyle`. To ensure portability to other PostgreSQL installations that might use non-default `DateStyle` settings, `DateStyle` should be set to `ISO` before using `COPY TO`. It is also a good idea to avoid dumping data with `IntervalStyle` set to `sql_standard`, because negative interval values might be misinterpreted by a server that has a different setting for `IntervalStyle`.

Input data is interpreted according to `ENCODING` option or the current client encoding, and output data is encoded in `ENCODING` or the current client encoding, even if the data does not pass through the client but is read from or written to a file directly by the server.

`COPY` stops operation at the first error. This should not lead to problems in the event of a `COPY TO`, but the target table will already have received earlier rows in a `COPY FROM`. These rows will not be visible or accessible, but they still occupy disk space. This might amount to a considerable amount of wasted disk space if the failure happened well into a large copy operation. You might wish to invoke `VACUUM` to recover the wasted space.

`FORCE_NULL` and `FORCE_NOT_NULL` can be used simultaneously on the same column. This results in converting quoted null strings to null values and unquoted null strings to empty strings.

## File Formats

### Text Format

When the text format is used, the data read or written is a text file with one line per table row. Columns in a row are separated by the delimiter character. The column values themselves are strings generated by the output function, or acceptable to the input function, of each attribute’s data type. The specified null string is used in place of columns that are null. `COPY FROM` will raise an error if any line of the input file contains more or fewer columns than are expected. If `OIDS` is specified, the OID is read or written as the first column, preceding the user data columns.

End of data can be represented by a single line containing just backslash-period (\.). An end-of-data marker is not necessary when reading from a file, since the end of file serves perfectly well; it is needed only when copying data to or from client applications using pre-3.0 client protocol.

Backslash characters (\) can be used in the `COPY` data to quote data characters that might otherwise be taken as row or column delimiters. In particular, the following characters must be preceded by a backslash if they appear as part of a column value: backslash itself, newline, carriage return, and the current delimiter character.

The specified null string is sent by `COPY TO` without adding any backslashes; conversely, `COPY FROM` matches the input against the null string before removing backslashes. Therefore, a null string such as `\N` cannot be confused with the actual data value `\N` (which would be represented as `\\N`).

The following special backslash sequences are recognized by `COPY FROM`:

<table>
<thead>
<tr>
<th>Sequence</th>
<th>Represents</th>
</tr>
</thead>
<tbody>
<tr>
<td>\b</td>
<td>Backspace (ASCII 8)</td>
</tr>
<tr>
<td>\f</td>
<td>Form feed (ASCII 12)</td>
</tr>
</tbody>
</table>
Presently, `COPY TO` will never emit an octal or hex-digits backslash sequence, but it does use the other sequences listed above for those control characters.

Any other backslashed character that is not mentioned in the above table will be taken to represent itself. However, beware of adding backslashes unnecessarily, since that might accidentally produce a string matching the end-of-data marker (`\.`) or the null string (`\n` by default). These strings will be recognized before any other backslash processing is done.

It is strongly recommended that applications generating `COPY` data convert data newlines and carriage returns to the `\n` and `\r` sequences respectively. At present it is possible to represent a data carriage return by a backslash and carriage return, and to represent a data newline by a backslash and newline. However, these representations might not be accepted in future releases. They are also highly vulnerable to corruption if the `COPY` file is transferred across different machines (for example, from Unix to Windows or vice versa).

`COPY TO` will terminate each row with a Unix-style newline ("\n"). Servers running on Microsoft Windows instead output carriage return/newline ("\r\n"), but only for `COPY` to a server file; for consistency across platforms, `COPY TO STDOUT` always sends "\n" regardless of server platform.

`COPY FROM` can handle lines ending with newlines, carriage returns, or carriage return/newlines. To reduce the risk of error due to un-backslashed newlines or carriage returns that were meant as data, `COPY FROM` will complain if the line endings in the input are not all alike.

### CSV Format

This format option is used for importing and exporting the Comma Separated Value (CSV) file format used by many other programs, such as spreadsheets. Instead of the escaping rules used by PostgreSQL’s standard text format, it produces and recognizes the common CSV escaping mechanism.

The values in each record are separated by the **delimiter** character. If the value contains the delimiter character, the **quote** character, the **null** string, a carriage return, or line feed character, then the whole value is prefixed and suffixed by the quote character, and any occurrence within the value of a quote character or the escape character is preceded by the escape character. You can also use **force_quote** to force quotes when outputting non-null values in specific columns.

The CSV format has no standard way to distinguish a null value from an empty string. PostgreSQL’s `COPY` handles this by quoting. A null is output as the null parameter string and is not quoted, while a non-null value matching the null parameter string is quoted. For example, with the default settings, a null is written as an unquoted empty string, while an empty string data value is written with double quotes ("""). Reading values follows similar rules. You can use **force_not_null** to prevent null input comparisons for specific columns. You can also use **force_null** to convert quoted null string data values to null.

Because backslash is not a special character in the CSV format, \., the end-of-data marker, could also appear as a data value. To avoid any misinterpretation, a \. data value appearing as a lone entry on
a line is automatically quoted on output, and on input, if quoted, is not interpreted as the end-of-data marker. If you are loading a file created by another application that has a single unquoted column and might have a value of \\\. you might need to quote that value in the input file.

**Note:** In CSV format, all characters are significant. A quoted value surrounded by white space, or any characters other than DELIMITER, will include those characters. This can cause errors if you import data from a system that pads CSV lines with white space out to some fixed width. If such a situation arises you might need to preprocess the CSV file to remove the trailing white space, before importing the data into PostgreSQL.

**Note:** CSV format will both recognize and produce CSV files with quoted values containing embedded carriage returns and line feeds. Thus the files are not strictly one line per table row like text-format files.

**Note:** Many programs produce strange and occasionally perverse CSV files, so the file format is more a convention than a standard. Thus you might encounter some files that cannot be imported using this mechanism, and COPY might produce files that other programs cannot process.

### Binary Format

The binary format option causes all data to be stored/read as binary format rather than as text. It is somewhat faster than the text and CSV formats, but a binary-format file is less portable across machine architectures and PostgreSQL versions. Also, the binary format is very data type specific; for example it will not work to output binary data from a smallint column and read it into an integer column, even though that would work fine in text format.

The binary file format consists of a file header, zero or more tuples containing the row data, and a file trailer. Headers and data are in network byte order.

**Note:** PostgreSQL releases before 7.4 used a different binary file format.

#### File Header

The file header consists of 15 bytes of fixed fields, followed by a variable-length header extension area. The fixed fields are:

**Signature**

11-byte sequence PGCOPY\n\377\r\n\0 — note that the zero byte is a required part of the signature. (The signature is designed to allow easy identification of files that have been munged by a non-8-bit-clean transfer. This signature will be changed by end-of-line-translation filters, dropped zero bytes, dropped high bits, or parity changes.)

**Flags field**

32-bit integer bit mask to denote important aspects of the file format. Bits are numbered from 0 (LSB) to 31 (MSB). Note that this field is stored in network byte order (most significant byte first), as are all the integer fields used in the file format. Bits 16-31 are reserved to denote critical
file format issues; a reader should abort if it finds an unexpected bit set in this range. Bits 0-15 are reserved to signal backwards-compatible format issues; a reader should simply ignore any unexpected bits set in this range. Currently only one flag bit is defined, and the rest must be zero:

Bit 16
if 1, OIDs are included in the data; if 0, not

Header extension area length
32-bit integer, length in bytes of remainder of header, not including self. Currently, this is zero, and the first tuple follows immediately. Future changes to the format might allow additional data to be present in the header. A reader should silently skip over any header extension data it does not know what to do with.

The header extension area is envisioned to contain a sequence of self-identifying chunks. The flags field is not intended to tell readers what is in the extension area. Specific design of header extension contents is left for a later release.

This design allows for both backwards-compatible header additions (add header extension chunks, or set low-order flag bits) and non-backwards-compatible changes (set high-order flag bits to signal such changes, and add supporting data to the extension area if needed).

Tuples
Each tuple begins with a 16-bit integer count of the number of fields in the tuple. (Presently, all tuples in a table will have the same count, but that might not always be true.) Then, repeated for each field in the tuple, there is a 32-bit length word followed by that many bytes of field data. (The length word does not include itself, and can be zero.) As a special case, -1 indicates a NULL field value. No value bytes follow in the NULL case.

There is no alignment padding or any other extra data between fields.

Presently, all data values in a binary-format file are assumed to be in binary format (format code one). It is anticipated that a future extension might add a header field that allows per-column format codes to be specified.

To determine the appropriate binary format for the actual tuple data you should consult the PostgreSQL source, in particular the \*send and \*recv functions for each column’s data type (typically these functions are found in the src/backend/utils/adt/ directory of the source distribution).

If OIDs are included in the file, the OID field immediately follows the field-count word. It is a normal field except that it’s not included in the field-count. In particular it has a length word — this will allow handling of 4-byte vs. 8-byte OIDs without too much pain, and will allow OIDs to be shown as null if that ever proves desirable.

File Trailer
The file trailer consists of a 16-bit integer word containing -1. This is easily distinguished from a tuple’s field-count word.

A reader should report an error if a field-count word is neither -1 nor the expected number of columns. This provides an extra check against somehow getting out of sync with the data.
Examples

The following example copies a table to the client using the vertical bar (|) as the field delimiter:

```
COPY country TO STDOUT (DELIMITER '|');
```

To copy data from a file into the `country` table:

```
COPY country FROM '/usr1/proj/bray/sql/country_data';
```

To copy into a file just the countries whose names start with 'A':

```
COPY (SELECT * FROM country WHERE country_name LIKE 'A%') TO '/usr1/proj/bray/sql/a_list_countries.copy';
```

To copy into a compressed file, you can pipe the output through an external compression program:

```
COPY country TO PROGRAM 'gzip > /usr1/proj/bray/sql/country_data.gz';
```

Here is a sample of data suitable for copying into a table from `STDIN`:

```
AF AFGHANISTAN
AL ALBANIA
DZ ALGERIA
ZM ZAMBIA
ZW ZIMBABWE
```

Note that the white space on each line is actually a tab character.

The following is the same data, output in binary format. The data is shown after filtering through the Unix utility `od -c`. The table has three columns; the first has type `char(2)`, the second has type `text`, and the third has type `integer`. All the rows have a null value in the third column.

```
0000000 P G C O P Y \n 377 \n \n \n \n \n \n 0000020 \n 0000040 F G H A N I S T A N 377 377 377 377 \n 0000060 \n 0000080 A L G E R I A 377 377 377 377 \n 0000100 A L B A N I A 377 377 377 377 \n 0000120 Z M Z A M B I A 377 377 377 377 \n 0000140 Z M Z I M B A B W E 377 377 377 377 377 377 377 377
```

Compatibility

There is no `COPY` statement in the SQL standard.

The following syntax was used before PostgreSQL version 9.0 and is still supported:

```
COPY table_name [ ( column_name [, ...] ) ]
    FROM { 'filename' | STDIN }
```
COPY

COPY { table_name ( ( column_name [, ...] ) | ( query ) ) }
TO { 'filename' | STDOUT }

Note that in this syntax, BINARY and CSV are treated as independent keywords, not as arguments of a FORMAT option.

The following syntax was used before PostgreSQL version 7.3 and is still supported:

COPY [ BINARY ] table_name [ WITH OIDS ]
FROM { 'filename' | STDIN }
[ [USING] DELIMITERS 'delimiter_character' ]
WITH NULL AS 'null_string'

COPY [ BINARY ] table_name [ WITH OIDS ]
TO { 'filename' | STDOUT }
[ [USING] DELIMITERS 'delimiter_character' ]
WITH NULL AS 'null_string'
CREATE ACCESS METHOD

Name
CREATE ACCESS METHOD — define a new access method

Synopsis
CREATE ACCESS METHOD name
    TYPE access_method_type
    HANDLER handler_function

Description
CREATE ACCESS METHOD creates a new access method.
The access method name must be unique within the database.
Only superusers can define new access methods.

Parameters

    name
    The name of the access method to be created.

    access_method_type
    This clause specifies the type of access method to define. Only INDEX is supported at present.

    handler_function
    handler_function is the name (possibly schema-qualified) of a previously registered function that represents the access method. The handler function must be declared to take a single argument of type internal, and its return type depends on the type of access method; for INDEX access methods, it must be index_am_handler. The C-level API that the handler function must implement varies depending on the type of access method. The index access method API is described in Chapter 59.

Examples
Create an index access method heptree with handler function heptree_handler:

CREATE ACCESS METHOD heptree TYPE INDEX HANDLER heptree_handler;

Compatibility
CREATE ACCESS METHOD is a PostgreSQL extension.
See Also

DROP ACCESS METHOD, CREATE OPERATOR CLASS, CREATE OPERATOR FAMILY
CREATE AGGREGATE

Name
CREATE AGGREGATE — define a new aggregate function

Synopsis
CREATE AGGREGATE name ( [ argmode ] [ argname ] arg_data_type [ , ... ] ) ( 
    SFUNC = sfunc,
    STYPE = state_data_type
    [ , SSPACE = state_data_size ]
    [ , FINALFUNC = ffunc ]
    [ , FINALFUNC_EXTRA ]
    [ , COMBINEFUNC = combinefunc ]
    [ , SERIALFUNC = serialfunc ]
    [ , DESERIALFUNC = deserialfunc ]
    [ , INITCOND = initial_condition ]
    [ , MSFUNC = mfunc ]
    [ , MINVFUNC = minvfunc ]
    [ , MSTYPE = mstate_data_type ]
    [ , MSSPACE = mstate_data_size ]
    [ , MFINALFUNC = mfunc ]
    [ , MFINALFUNC_EXTRA ]
    [ , INITCOND = initial_condition ]
    [ , PARALLEL = { SAFE | RESTRICTED | UNSAFE } ]
) 

or the old syntax

CREATE AGGREGATE name ( [ argmode ] [ argname ] arg_data_type [ , ... ] ) ( 
    SFUNC = sfunc,
    STYPE = state_data_type
    [ , SSPACE = state_data_size ]
    [ , FINALFUNC = ffunc ]
    [ , FINALFUNC_EXTRA ]
    [ , INITCOND = initial_condition ]
    [ , PARALLEL = { SAFE | RESTRICTED | UNSAFE } ]
    [ , HYPOTHETICAL ]
) 

or the old syntax

CREATE AGGREGATE name ( 
    BASETYPE = base_type,
    SFUNC = sfunc,
    STYPE = state_data_type
    [ , SSPACE = state_data_size ]
    [ , FINALFUNC = ffunc ]
    [ , FINALFUNC_EXTRA ]
    [ , COMBINEFUNC = combinefunc ]
    [ , SERIALFUNC = serialfunc ]
    [ , DESERIALFUNC = deserialfunc ]
    [ , INITCOND = initial_condition ]
    [ , MSFUNC = mfunc ]
CREATE AGGREGATE

[ , MINVFUNC = minvfunc ]
[ , MSTYPE = mstate_data_type ]
[ , MSSPACE = mstate_data_size ]
[ , MFINALFUNC = mffunc ]
[ , MFINALFUNC_EXTRA ]
[ , MINITCOND = minitial_condition ]
[ , SORTOP = sort_operator ]
}

Description

CREATE AGGREGATE defines a new aggregate function. Some basic and commonly-used aggregate functions are included with the distribution; they are documented in Section 9.20. If one defines new types or needs an aggregate function not already provided, then CREATE AGGREGATE can be used to provide the desired features.

If a schema name is given (for example, CREATE AGGREGATE myschema.myagg ...) then the aggregate function is created in the specified schema. Otherwise it is created in the current schema.

An aggregate function is identified by its name and input data type(s). Two aggregates in the same schema can have the same name if they operate on different input types. The name and input data type(s) of an aggregate must also be distinct from the name and input data type(s) of every ordinary function in the same schema. This behavior is identical to overloading of ordinary function names (see CREATE FUNCTION).

A simple aggregate function is made from one or two ordinary functions: a state transition function sfunc, and an optional final calculation function ffunc. These are used as follows:

\[
\text{sfunc}( \text{internal-state}, \text{next-data-values} ) \rightarrow \text{next-internal-state} \\
\text{ffunc}( \text{internal-state} ) \rightarrow \text{aggregate-value}
\]

PostgreSQL creates a temporary variable of data type stype to hold the current internal state of the aggregate. At each input row, the aggregate argument value(s) are calculated and the state transition function is invoked with the current state value and the new argument value(s) to calculate a new internal state value. After all the rows have been processed, the final function is invoked once to calculate the aggregate’s return value. If there is no final function then the ending state value is returned as-is.

An aggregate function can provide an initial condition, that is, an initial value for the internal state value. This is specified and stored in the database as a value of type text, but it must be a valid external representation of a constant of the state value data type. If it is not supplied then the state value starts out null.

If the state transition function is declared “strict”, then it cannot be called with null inputs. With such a transition function, aggregate execution behaves as follows. Rows with any null input values are ignored (the function is not called and the previous state value is retained). If the initial state value is null, then at the first row with all-nonnull input values, the first argument value replaces the state value, and the transition function is invoked at each subsequent row with all-nonnull input values. This is handy for implementing aggregates like \text{max}. Note that this behavior is only available when stype is the same as the first arg_data_type. When these types are different, you must supply a nonnull initial condition or use a nonstrict transition function.

If the state transition function is not strict, then it will be called unconditionally at each input row, and must deal with null inputs and null state values for itself. This allows the aggregate author to have full control over the aggregate’s handling of null values.
CREATE AGGREGATE

If the final function is declared "strict", then it will not be called when the ending state value is null; instead a null result will be returned automatically. (Of course this is just the normal behavior of strict functions.) In any case the final function has the option of returning a null value. For example, the final function for \texttt{avg} returns null when it sees there were zero input rows.

Sometimes it is useful to declare the final function as taking not just the state value, but extra parameters corresponding to the aggregate’s input values. The main reason for doing this is if the final function is polymorphic and the state value’s data type would be inadequate to pin down the result type. These extra parameters are always passed as NULL (and so the final function must not be strict when the \texttt{FINALFUNC\_EXTRA} option is used), but nonetheless they are valid parameters. The final function could for example make use of \texttt{get\_fn\_expr\_argtype} to identify the actual argument type in the current call.

An aggregate can optionally support \textit{moving-aggregate mode}, as described in Section 36.10.1. This requires specifying the \texttt{MSFUNC}, \texttt{MINVFUNC}, and \texttt{MSTYPE} parameters, and optionally the \texttt{MSPACE}, \texttt{MFINALFUNC}, \texttt{MFINALFUNC\_EXTRA}, and \texttt{MINITCOND} parameters. Except for \texttt{MINVFUNC}, these parameters work like the corresponding simple-aggregate parameters without \texttt{M}; they define a separate implementation of the aggregate that includes an inverse transition function.

The syntax with \texttt{ORDER\ BY} in the parameter list creates a special type of aggregate called an \textit{ordered-set aggregate}; or if \texttt{HYPOTHETICAL} is specified, then a \textit{hypothetical-set aggregate} is created. These aggregates operate over groups of sorted values in order-dependent ways, so that specification of an input sort order is an essential part of a call. Also, they can have \textit{direct} arguments, which are arguments that are evaluated only once per aggregation rather than once per input row. Hypothetical-set aggregates are a subclass of ordered-set aggregates in which some of the direct arguments are required to match, in number and data types, the aggregated argument columns. This allows the values of those direct arguments to be added to the collection of aggregate-input rows as an additional “hypothetical” row.

An aggregate can optionally support \textit{partial aggregation}, as described in Section 36.10.4. This requires specifying the \texttt{COMBINEFUNC} parameter. If the \texttt{state\_data\_type} is \texttt{internal}, it’s usually also appropriate to provide the \texttt{SERIALFUNC} and \texttt{DESERIALFUNC} parameters so that parallel aggregation is possible. Note that the aggregate must also be marked \texttt{PARALLEL\ SAFE} to enable parallel aggregation.

An aggregate can behave like \texttt{MIN} or \texttt{MAX} can sometimes be optimized by looking into an index instead of scanning every input row. If this aggregate can be so optimized, indicate it by specifying a \textit{sort operator}. The basic requirement is that the aggregate must yield the first element in the sort ordering induced by the operator; in other words:

\begin{verbatim}
SELECT agg(col) FROM tab;
\end{verbatim}

must be equivalent to:

\begin{verbatim}
SELECT col FROM tab ORDER BY col USING sortop LIMIT 1;
\end{verbatim}

Further assumptions are that the aggregate ignores null inputs, and that it delivers a null result if and only if there were no non-null inputs. Ordinarily, a data type’s \texttt{<} operator is the proper sort operator for \texttt{MIN}, and \texttt{>} is the proper sort operator for \texttt{MAX}. Note that the optimization will never actually take effect unless the specified operator is the “less than” or “greater than” strategy member of a B-tree index operator class.

To be able to create an aggregate function, you must have \texttt{USAGE} privilege on the argument types, the state type(s), and the return type, as well as \texttt{EXECUTE} privilege on the supporting functions.
CREATE AGGREGATE

Parameters

name
The name (optionally schema-qualified) of the aggregate function to create.

argmode
The mode of an argument: IN or VARIADIC. (Aggregate functions do not support OUT arguments.) If omitted, the default is IN. Only the last argument can be marked VARIADIC.

argname
The name of an argument. This is currently only useful for documentation purposes. If omitted, the argument has no name.

arg_data_type
An input data type on which this aggregate function operates. To create a zero-argument aggregate function, write * in place of the list of argument specifications. (An example of such an aggregate is \texttt{count(*)}.)

base_type
In the old syntax for CREATE AGGREGATE, the input data type is specified by a basetype parameter rather than being written next to the aggregate name. Note that this syntax allows only one input parameter. To define a zero-argument aggregate function with this syntax, specify the basetype as "ANY" (not *). Ordered-set aggregates cannot be defined with the old syntax.

sfunc
The name of the state transition function to be called for each input row. For a normal \(N\)-argument aggregate function, the \textit{sfunc} must take \(N+1\) arguments, the first being of type \textit{state_data_type} and the rest matching the declared input data type(s) of the aggregate. The function must return a value of type \textit{state_data_type}. This function takes the current state value and the current input data value(s), and returns the next state value.

For ordered-set (including hypothetical-set) aggregates, the state transition function receives only the current state value and the aggregated arguments, not the direct arguments. Otherwise it is the same.

\textit{state_data_type}
The data type for the aggregate’s state value.

\textit{state_data_size}
The approximate average size (in bytes) of the aggregate’s state value. If this parameter is omitted or is zero, a default estimate is used based on the \textit{state_data_type}. The planner uses this value to estimate the memory required for a grouped aggregate query. The planner will consider using hash aggregation for such a query only if the hash table is estimated to fit in \texttt{work_mem}; therefore, large values of this parameter discourage use of hash aggregation.

ffunc
The name of the final function called to compute the aggregate’s result after all input rows have been traversed. For a normal aggregate, this function must take a single argument of type \textit{state_data_type}. The return data type of the aggregate is defined as the return type of this function. If \textit{ffunc} is not specified, then the ending state value is used as the aggregate’s result, and the return type is \textit{state_data_type}.

For ordered-set (including hypothetical-set) aggregates, the final function receives not only the final state value, but also the values of all the direct arguments.
CREATE AGGREGATE

If `FINALFUNC_EXTRA` is specified, then in addition to the final state value and any direct arguments, the final function receives extra NULL values corresponding to the aggregate’s regular (aggregated) arguments. This is mainly useful to allow correct resolution of the aggregate result type when a polymorphic aggregate is being defined.

**combinefunc**

The `combinefunc` function may optionally be specified to allow the aggregate function to support partial aggregation. If provided, the `combinefunc` must combine two `state_data_type` values, each containing the result of aggregation over some subset of the input values, to produce a new `state_data_type` that represents the result of aggregating over both sets of inputs. This function can be thought of as an `sfunc`, where instead of acting upon an individual input row and adding it to the running aggregate state, it adds another aggregate state to the running state.

The `combinefunc` must be declared as taking two arguments of the `state_data_type` and returning a value of the `state_data_type`. Optionally this function may be “strict”. In this case the function will not be called when either of the input states are null; the other state will be taken as the correct result.

For aggregate functions whose `state_data_type` is `internal`, the `combinefunc` must not be strict. In this case the `combinefunc` must ensure that null states are handled correctly and that the state being returned is properly stored in the aggregate memory context.

**serialfunc**

An aggregate function whose `state_data_type` is `internal` can participate in parallel aggregation only if it has a `serialfunc` function, which must serialize the aggregate state into a `bytea` value for transmission to another process. This function must take a single argument of type `internal` and return type `bytea`. A corresponding `deserialfunc` is also required.

**deserialfunc**

Deserializes a previously serialized aggregate state back into `state_data_type`. This function must take two arguments of types `bytea` and `internal`, and produce a result of type `internal`. (Note: the second, `internal` argument is unused, but is required for type safety reasons.)

**initial_condition**

The initial setting for the state value. This must be a string constant in the form accepted for the data type `state_data_type`. If not specified, the state value starts out null.

**msfunc**

The name of the forward state transition function to be called for each input row in moving-aggregate mode. This is exactly like the regular transition function, except that its first argument and result are of type `mstate_data_type`, which might be different from `state_data_type`.

**minvfunc**

The name of the inverse state transition function to be used in moving-aggregate mode. This function has the same argument and result types as `msfunc`, but it is used to remove a value from the current aggregate state, rather than add a value to it. The inverse transition function must have the same strictness attribute as the forward state transition function.

**mstate_data_type**

The data type for the aggregate’s state value, when using moving-aggregate mode.

**mstate_data_size**

The approximate average size (in bytes) of the aggregate’s state value, when using moving-aggregate mode. This works the same as `state_data_size`. 
The name of the final function called to compute the aggregate’s result after all input rows have been traversed, when using moving-aggregate mode. This works the same as `ffunc`, except that its first argument’s type is `mstate_data_type` and extra dummy arguments are specified by writing `MFINALFUNC_EXTRA`. The aggregate result type determined by `mffunc` or `mstate_data_type` must match that determined by the aggregate’s regular implementation.

The initial setting for the state value, when using moving-aggregate mode. This works the same as `initial_condition`.

The associated sort operator for a `MIN`- or `MAX`-like aggregate. This is just an operator name (possibly schema-qualified). The operator is assumed to have the same input data types as the aggregate (which must be a single-argument normal aggregate).

The meanings of `PARALLEL SAFE`, `PARALLEL RESTRICTED`, and `PARALLEL UNSAFE` are the same as for `CREATE FUNCTION`. An aggregate will not be considered for parallelization if it is marked `PARALLEL UNSAFE` (which is the default!) or `PARALLEL RESTRICTED`. Note that the parallel-safety markings of the aggregate’s support functions are not consulted by the planner, only the marking of the aggregate itself.

For ordered-set aggregates only, this flag specifies that the aggregate arguments are to be processed according to the requirements for hypothetical-set aggregates: that is, the last few direct arguments must match the data types of the aggregated (`WITHIN GROUP`) arguments. The `HYPOTHETICAL` flag has no effect on run-time behavior, only on parse-time resolution of the data types and collations of the aggregate’s arguments.

The parameters of `CREATE AGGREGATE` can be written in any order, not just the order illustrated above.

In parameters that specify support function names, you can write a schema name if needed, for example `SFUNC = public.sum`. Do not write argument types there, however — the argument types of the support functions are determined from other parameters.

If an aggregate supports moving-aggregate mode, it will improve calculation efficiency when the aggregate is used as a window function for a window with moving frame start (that is, a frame start mode other than `UNBOUNDED PRECEDING`). Conceptually, the forward transition function adds input values to the aggregate’s state when they enter the window frame from the bottom, and the inverse transition function removes them again when they leave the frame at the top. So, when values are removed, they are always removed in the same order they were added. Whenever the inverse transition function is invoked, it will thus receive the earliest added but not yet removed argument value(s). The inverse transition function can assume that at least one row will remain in the current state after it removes the oldest row. (When this would not be the case, the window function mechanism simply starts a fresh aggregation, rather than using the inverse transition function.)

The forward transition function for moving-aggregate mode is not allowed to return `NULL` as the new state value. If the inverse transition function returns `NULL`, this is taken as an indication that the inverse function cannot reverse the state calculation for this particular input, and so the aggregate
calculation will be redone from scratch for the current frame starting position. This convention al-
allows moving-aggregate mode to be used in situations where there are some infrequent cases that are
impractical to reverse out of the running state value.

If no moving-aggregate implementation is supplied, the aggregate can still be used with moving
frames, but PostgreSQL will recompute the whole aggregation whenever the start of the frame moves.
Note that whether or not the aggregate supports moving-aggregate mode, PostgreSQL can handle a
moving frame end without recalculation; this is done by continuing to add new values to the aggre-
gate’s state. It is assumed that the final function does not damage the aggregate’s state value, so that
the aggregation can be continued even after an aggregate result value has been obtained for one set of
frame boundaries.

The syntax for ordered-set aggregates allows VARIADIC to be specified for both the last direct pa-
rameter and the last aggregated (WITHIN GROUP) parameter. However, the current implementation
restricts use of VARIADIC in two ways. First, ordered-set aggregates can only use VARIADIC "any",
not other variadic array types. Second, if the last direct parameter is VARIADIC "any", then there
can be only one aggregated parameter and it must also be VARIADIC "any". (In the representation
used in the system catalogs, these two parameters are merged into a single VARIADIC "any" item,
since pg_proc cannot represent functions with more than one VARIADIC parameter.) If the aggregate
is a hypothetical-set aggregate, the direct arguments that match the VARIADIC "any" parameter are
the hypothetical ones; any preceding parameters represent additional direct arguments that are not
constrained to match the aggregated arguments.

Currently, ordered-set aggregates do not need to support moving-aggregate mode, since they cannot
be used as window functions.

Partial (including parallel) aggregation is currently not supported for ordered-set aggregates. Also,
it will never be used for aggregate calls that include DISTINCT or ORDER BY clauses, since those
semantics cannot be supported during partial aggregation.

Examples

See Section 36.10.

Compatibility

CREATE AGGREGATE is a PostgreSQL language extension. The SQL standard does not provide for
user-defined aggregate functions.

See Also

ALTER AGGREGATE, DROP AGGREGATE
CREATE CAST

Name
CREATE CAST — define a new cast

Synopsis
CREATE CAST (source_type AS target_type)
    WITH FUNCTION function_name (argument_type [, ...])
    [ AS ASSIGNMENT | AS IMPLICIT ]

CREATE CAST (source_type AS target_type)
    WITHOUT FUNCTION
    [ AS ASSIGNMENT | AS IMPLICIT ]

CREATE CAST (source_type AS target_type)
    WITH INOUT
    [ AS ASSIGNMENT | AS IMPLICIT ]

Description
CREATE CAST defines a new cast. A cast specifies how to perform a conversion between two data types. For example,

```
SELECT CAST(42 AS float8);
```

converts the integer constant 42 to type float8 by invoking a previously specified function, in this case float8(int4). (If no suitable cast has been defined, the conversion fails.)

Two types can be binary coercible, which means that the conversion can be performed “for free” without invoking any function. This requires that corresponding values use the same internal representation. For instance, the types text and varchar are binary coercible both ways. Binary coercibility is not necessarily a symmetric relationship. For example, the cast from xml to text can be performed for free in the present implementation, but the reverse direction requires a function that performs at least a syntax check. (Two types that are binary coercible both ways are also referred to as binary compatible.)

You can define a cast as an I/O conversion cast by using the WITH INOUT syntax. An I/O conversion cast is performed by invoking the output function of the source data type, and passing the resulting string to the input function of the target data type. In many common cases, this feature avoids the need to write a separate cast function for conversion. An I/O conversion cast acts the same as a regular function-based cast; only the implementation is different.

By default, a cast can be invoked only by an explicit cast request, that is an explicit CAST(x AS typename) or x::typename construct.

If the cast is marked AS ASSIGNMENT then it can be invoked implicitly when assigning a value to a column of the target data type. For example, supposing that foo.f1 is a column of type text, then:

```
INSERT INTO foo (f1) VALUES (42);
```
CREATE CAST

will be allowed if the cast from type integer to type text is marked AS ASSIGNMENT, otherwise not. (We generally use the term assignment cast to describe this kind of cast.)

If the cast is marked AS IMPLICIT then it can be invoked implicitly in any context, whether assignment or internally in an expression. (We generally use the term implicit cast to describe this kind of cast.) For example, consider this query:

```
SELECT 2 + 4.0;
```

The parser initially marks the constants as being of type integer and numeric respectively. There is no integer + numeric operator in the system catalogs, but there is a numeric + numeric operator. The query will therefore succeed if a cast from integer to numeric is available and is marked AS IMPLICIT — which in fact it is. The parser will apply the implicit cast and resolve the query as if it had been written

```
SELECT CAST ( 2 AS numeric ) + 4.0;
```

Now, the catalogs also provide a cast from numeric to integer. If that cast were marked AS IMPLICIT — which it is not — then the parser would be faced with choosing between the above interpretation and the alternative of casting the numeric constant to integer and applying the integer + integer operator. Lacking any knowledge of which choice to prefer, it would give up and declare the query ambiguous. The fact that only one of the two casts is implicit is the way in which we teach the parser to prefer resolution of a mixed numeric-and-integer expression as numeric; there is no built-in knowledge about that.

It is wise to be conservative about marking casts as implicit. An overabundance of implicit casting paths can cause PostgreSQL to choose surprising interpretations of commands, or to be unable to resolve commands at all because there are multiple possible interpretations. A good rule of thumb is to make a cast implicitly invokable only for information-preserving transformations between types in the same general type category. For example, the cast from int2 to int4 can reasonably be implicit, but the cast from float8 to int4 should probably be assignment-only. Cross-type-category casts, such as text to int4, are best made explicit-only.

**Note:** Sometimes it is necessary for usability or standards-compliance reasons to provide multiple implicit casts among a set of types, resulting in ambiguity that cannot be avoided as above. The parser has a fallback heuristic based on type categories and preferred types that can help to provide desired behavior in such cases. See CREATE TYPE for more information.

To be able to create a cast, you must own the source or the target data type and have USAGE privilege on the other type. To create a binary-coercible cast, you must be superuser. (This restriction is made because an erroneous binary-coercible cast conversion can easily crash the server.)

**Parameters**

*source_type*

The name of the source data type of the cast.

*target_type*

The name of the target data type of the cast.
CREATE CAST

function_name(argument_type [, ...])

The function used to perform the cast. The function name can be schema-qualified. If it is not, the function will be looked up in the schema search path. The function’s result data type must match the target type of the cast. Its arguments are discussed below.

WITHOUT FUNCTION

Indicates that the source type is binary-coercible to the target type, so no function is required to perform the cast.

WITH INOUT

Indicates that the cast is an I/O conversion cast, performed by invoking the output function of the source data type, and passing the resulting string to the input function of the target data type.

AS ASSIGNMENT

Indicates that the cast can be invoked implicitly in assignment contexts.

AS IMPLICIT

Indicates that the cast can be invoked implicitly in any context.

Cast implementation functions can have one to three arguments. The first argument type must be identical to or binary-coercible from the cast’s source type. The second argument, if present, must be type integer; it receives the type modifier associated with the destination type, or -1 if there is none. The third argument, if present, must be type boolean; it receives true if the cast is an explicit cast, false otherwise. (Bizarrely, the SQL standard demands different behaviors for explicit and implicit casts in some cases. This argument is supplied for functions that must implement such casts. It is not recommended that you design your own data types so that this matters.)

The return type of a cast function must be identical to or binary-coercible to the cast’s target type.

Ordinarily a cast must have different source and target data types. However, it is allowed to declare a cast with identical source and target types if it has a cast implementation function with more than one argument. This is used to represent type-specific length coercion functions in the system catalogs. The named function is used to coerce a value of the type to the type modifier value given by its second argument.

When a cast has different source and target types and a function that takes more than one argument, it supports converting from one type to another and applying a length coercion in a single step. When no such entry is available, coercion to a type that uses a type modifier involves two cast steps, one to convert between data types and a second to apply the modifier.

A cast to or from a domain type currently has no effect. Casting to or from a domain uses the casts associated with its underlying type.

Notes

Use DROP CAST to remove user-defined casts.

Remember that if you want to be able to convert types both ways you need to declare casts both ways explicitly.

It is normally not necessary to create casts between user-defined types and the standard string types (text, varchar, and char(n), as well as user-defined types that are defined to be in the string category). PostgreSQL provides automatic I/O conversion casts for that. The automatic casts to string types are treated as assignment casts, while the automatic casts from string types are explicit-only. You can override this behavior by declaring your own cast to replace an automatic cast, but usually
the only reason to do so is if you want the conversion to be more easily invokable than the standard
assignment-only or explicit-only setting. Another possible reason is that you want the conversion to
behave differently from the type’s I/O function; but that is sufficiently surprising that you should think
twice about whether it’s a good idea. (A small number of the built-in types do indeed have different
behaviors for conversions, mostly because of requirements of the SQL standard.)

While not required, it is recommended that you continue to follow this old convention of naming
cast implementation functions after the target data type. Many users are used to being able to cast
data types using a function-style notation, that is typename(x). This notation is in fact nothing more
nor less than a call of the cast implementation function; it is not specially treated as a cast. If your
conversion functions are not named to support this convention then you will have surprised users.
Since PostgreSQL allows overloading of the same function name with different argument types, there
is no difficulty in having multiple conversion functions from different types that all use the target
type’s name.

**Note:** Actually the preceding paragraph is an oversimplification: there are two cases in which a
function-call construct will be treated as a cast request without having matched it to an actual
function. If a function call name(x) does not exactly match any existing function, but name is the
name of a data type and pg_cast provides a binary-coercible cast to this type from the type
of x, then the call will be construed as a binary-coercible cast. This exception is made so that
binary-coercible casts can be invoked using functional syntax, even though they lack any function.
Likewise, if there is no pg_cast entry but the cast would be to or from a string type, the call will
be construed as an I/O conversion cast. This exception allows I/O conversion casts to be invoked
using functional syntax.

**Note:** There is also an exception to the exception: I/O conversion casts from composite types
to string types cannot be invoked using functional syntax, but must be written in explicit cast
syntax (either CAST or :: notation). This exception was added because after the introduction of
automatically-provided I/O conversion casts, it was found too easy to accidentally invoke such a
cast when a function or column reference was intended.

**Examples**

To create an assignment cast from type bigint to type int4 using the function int4(bigint):

```
CREATE CAST (bigint AS int4) WITH FUNCTION int4(bigint) AS ASSIGNMENT;
```

(This cast is already predefined in the system.)

**Compatibility**

The `CREATE CAST` command conforms to the SQL standard, except that SQL does not make provi-
sions for binary-coercible types or extra arguments to implementation functions. AS IMPLICIT is a
PostgreSQL extension, too.
See Also

CREATE FUNCTION, CREATE TYPE, DROP CAST
CREATE COLLATION

Name
CREATE COLLATION — define a new collation

Synopsis
CREATE COLLATION name (  
    [ LOCALE = locale, ]  
    [ LC_COLLATE = lc_collate, ]  
    [ LC_CTYPE = lc_ctype ]  
  )  
CREATE COLLATION name FROM existing_collation

Description
CREATE COLLATION defines a new collation using the specified operating system locale settings, or by copying an existing collation.

To be able to create a collation, you must have CREATE privilege on the destination schema.

Parameters

name
   The name of the collation. The collation name can be schema-qualified. If it is not, the collation is defined in the current schema. The collation name must be unique within that schema. (The system catalogs can contain collations with the same name for other encodings, but these are ignored if the database encoding does not match.)

locale
   This is a shortcut for setting LC_COLLATE and LC_CTYPE at once. If you specify this, you cannot specify either of those parameters.

lc_collate
   Use the specified operating system locale for the LC_COLLATE locale category. The locale must be applicable to the current database encoding. (See CREATE DATABASE for the precise rules.)

lc_ctype
   Use the specified operating system locale for the LC_CTYPE locale category. The locale must be applicable to the current database encoding. (See CREATE DATABASE for the precise rules.)

existing_collation
   The name of an existing collation to copy. The new collation will have the same properties as the existing one, but it will be an independent object.
CREATE COLLATION

Notes
Use `DROP COLLATION` to remove user-defined collations.
See Section 23.2 for more information about collation support in PostgreSQL.

Examples
To create a collation from the operating system locale `fr_FR.utf8` (assuming the current database encoding is `UTF8`):

```sql
CREATE COLLATION french (LOCALE = 'fr_FR.utf8');
```

To create a collation from an existing collation:

```sql
CREATE COLLATION german FROM "de_DE";
```

This can be convenient to be able to use operating-system-independent collation names in applications.

Compatibility
There is a `CREATE COLLATION` statement in the SQL standard, but it is limited to copying an existing collation. The syntax to create a new collation is a PostgreSQL extension.

See Also
`ALTER COLLATION`, `DROP COLLATION`
CREATE CONVERSION

Name

CREATE CONVERSION — define a new encoding conversion

Synopsis

CREATE [ DEFAULT ] CONVERSION name
    FOR source_encoding TO dest_encoding FROM function_name

Description

CREATE CONVERSION defines a new conversion between character set encodings. Also, conversions
that are marked DEFAULT can be used for automatic encoding conversion between client and server.
For this purpose, two conversions, from encoding A to B and from encoding B to A, must be defined.
To be able to create a conversion, you must have EXECUTE privilege on the function and CREATE
privilege on the destination schema.

Parameters

DEFAULT

The DEFAULT clause indicates that this conversion is the default for this particular source to
destination encoding. There should be only one default encoding in a schema for the encoding
pair.

name

The name of the conversion. The conversion name can be schema-qualified. If it is not, the con-
version is defined in the current schema. The conversion name must be unique within a schema.

source_encoding

The source encoding name.

dest_encoding

The destination encoding name.

function_name

The function used to perform the conversion. The function name can be schema-qualified. If it is
not, the function will be looked up in the path.

The function must have the following signature:

conv_proc(
    integer, -- source encoding ID
    integer, -- destination encoding ID
    cstring, -- source string (null terminated C string)
    internal, -- destination (fill with a null terminated C string)
    integer -- source string length
) RETURNS void;
CREATE CONVERSION

Notes

Use DROP CONVERSION to remove user-defined conversions.
The privileges required to create a conversion might be changed in a future release.

Examples

To create a conversion from encoding UTF8 to LATIN1 using myfunc:

```sql
CREATE CONVERSION myconv FOR 'UTF8' TO 'LATIN1' FROM myfunc;
```

Compatibility

CREATE CONVERSION is a PostgreSQL extension. There is no CREATE CONVERSION statement in the SQL standard, but a CREATE TRANSLATION statement that is very similar in purpose and syntax.

See Also

ALTER CONVERSION, CREATE FUNCTION, DROP CONVERSION
CREATE DATABASE

Name

CREATE DATABASE — create a new database

Synopsis

CREATE DATABASE name
  [ [ WITH ] [ OWNER [=] user_name ]
  [ TEMPLATE [=] template ]
  [ ENCODING [=] encoding ]
  [ LC_COLLATE [=] lc_collate ]
  [ LC_CTYPE [=] lc_ctype ]
  [ TABLESPACE [=] tablespace_name ]
  [ ALLOW_CONNECTIONS [=] allowconn ]
  [ CONNECTION LIMIT [=] connlimit ]
  [ IS_TEMPLATE [=] istemplate ] ]

Description

CREATE DATABASE creates a new PostgreSQL database.

To create a database, you must be a superuser or have the special CREATEDB privilege. See CREATE USER.

By default, the new database will be created by cloning the standard system database template1. A different template can be specified by writing TEMPLATE name. In particular, by writing TEMPLATE template0, you can create a virgin database containing only the standard objects predefined by your version of PostgreSQL. This is useful if you wish to avoid copying any installation-local objects that might have been added to template1.

Parameters

name

The name of a database to create.

user_name

The role name of the user who will own the new database, or DEFAULT to use the default (namely, the user executing the command). To create a database owned by another role, you must be a direct or indirect member of that role, or be a superuser.

template

The name of the template from which to create the new database, or DEFAULT to use the default template (template1).

encoding

Character set encoding to use in the new database. Specify a string constant (e.g., ’SQL_ASCII’), or an integer encoding number, or DEFAULT to use the default encoding.
(namely, the encoding of the template database). The character sets supported by the PostgreSQL server are described in Section 23.3.1. See below for additional restrictions.

**lc_collate**

Collation order (`LC_COLLATE`) to use in the new database. This affects the sort order applied to strings, e.g. in queries with `ORDER BY`, as well as the order used in indexes on text columns. The default is to use the collation order of the template database. See below for additional restrictions.

**lc_ctype**

Character classification (`LC_CTYPE`) to use in the new database. This affects the categorization of characters, e.g. lower, upper and digit. The default is to use the character classification of the template database. See below for additional restrictions.

**tablespace_name**

The name of the tablespace that will be associated with the new database, or `DEFAULT` to use the template database’s tablespace. This tablespace will be the default tablespace used for objects created in this database. See `CREATE TABLESPACE` for more information.

**allowconn**

If false then no one can connect to this database. The default is true, allowing connections (except as restricted by other mechanisms, such as `GRANT/REVOKE CONNECT`).

**connlimit**

How many concurrent connections can be made to this database. -1 (the default) means no limit.

**istemplate**

If true, then this database can be cloned by any user with `CREATEDB` privileges; if false (the default), then only superusers or the owner of the database can clone it.

Optional parameters can be written in any order, not only the order illustrated above.

**Notes**

`CREATE DATABASE` cannot be executed inside a transaction block.

Errors along the line of “could not initialize database directory” are most likely related to insufficient permissions on the data directory, a full disk, or other file system problems.

Use `DROP DATABASE` to remove a database.

The program createdb is a wrapper program around this command, provided for convenience.

Database-level configuration parameters (set via `ALTER DATABASE`) are not copied from the template database.

Although it is possible to copy a database other than `template1` by specifying its name as the template, this is not (yet) intended as a general-purpose “COPY DATABASE” facility. The principal limitation is that no other sessions can be connected to the template database while it is being copied. `CREATE DATABASE` will fail if any other connection exists when it starts; otherwise, new connections to the template database are locked out until `CREATE DATABASE` completes. See Section 22.3 for more information.

The character set encoding specified for the new database must be compatible with the chosen locale settings (`LC_COLLATE` and `LC_CTYPE`). If the locale is `C` (or equivalently `POSIX`), then all encodings are allowed, but for other locale settings there is only one encoding that will work properly. (On Windows, however, UTF-8 encoding can be used with any locale.) `CREATE DATABASE` will allow supe-
CREATE DATABASE

users to specify SQL_ASCII encoding regardless of the locale settings, but this choice is deprecated and may result in misbehavior of character-string functions if data that is not encoding-compatible with the locale is stored in the database.

The encoding and locale settings must match those of the template database, except when template0 is used as template. This is because other databases might contain data that does not match the specified encoding, or might contain indexes whose sort ordering is affected by LC_COLLATE and LC_CTYPE. Copying such data would result in a database that is corrupt according to the new settings. template0, however, is known to not contain any data or indexes that would be affected.

The CONNECTION LIMIT option is only enforced approximately; if two new sessions start at about the same time when just one connection “slot” remains for the database, it is possible that both will fail. Also, the limit is not enforced against superusers or background worker processes.

Examples

To create a new database:

```
CREATE DATABASE lusiadas;
```

To create a database sales owned by user salesapp with a default tablespace of salesspace:

```
CREATE DATABASE sales OWNER salesapp TABLESPACE salesspace;
```

To create a database music which supports the ISO-8859-1 character set:

```
CREATE DATABASE music ENCODING 'LATIN1' TEMPLATE template0;
```

In this example, the TEMPLATE template0 clause would only be required if template1's encoding is not ISO-8859-1. Note that changing encoding might require selecting new LC_COLLATE and LC_CTYPE settings as well.

Compatibility

There is no CREATE DATABASE statement in the SQL standard. Databases are equivalent to catalogs, whose creation is implementation-defined.

See Also

ALTER DATABASE, DROP DATABASE
CREATE DOMAIN

**Name**
CREATE DOMAIN — define a new domain

**Synopsis**
CREATE DOMAIN name [ AS ] data_type  
[ COLLATE collation ]  
[ DEFAULT expression ]  
[ constraint [ ... ] ]

where constraint is:

[ CONSTRAINT constraint_name ]  
{ NOT NULL | NULL | CHECK (expression) }

**Description**
CREATE DOMAIN creates a new domain. A domain is essentially a data type with optional constraints (restrictions on the allowed set of values). The user who defines a domain becomes its owner.

If a schema name is given (for example, CREATE DOMAIN myschema.mydomain ...) then the domain is created in the specified schema. Otherwise it is created in the current schema. The domain name must be unique among the types and domains existing in its schema.

Domains are useful for abstracting common constraints on fields into a single location for maintenance. For example, several tables might contain email address columns, all requiring the same CHECK constraint to verify the address syntax. Define a domain rather than setting up each table’s constraint individually.

To be able to create a domain, you must have USAGE privilege on the underlying type.

**Parameters**

*name*

The name (optionally schema-qualified) of a domain to be created.

*data_type*

The underlying data type of the domain. This can include array specifiers.

*collation*

An optional collation for the domain. If no collation is specified, the underlying data type’s default collation is used. The underlying type must be collatable if COLLATE is specified.

DEFAULT *expression*

The DEFAULT clause specifies a default value for columns of the domain data type. The value is any variable-free expression (but subqueries are not allowed). The data type of the default
expression must match the data type of the domain. If no default value is specified, then the
default value is the null value.

The default expression will be used in any insert operation that does not specify a value for the
column. If a default value is defined for a particular column, it overrides any default associated
with the domain. In turn, the domain default overrides any default value associated with the
underlying data type.

CONSTRAINT `constraint_name`

An optional name for a constraint. If not specified, the system generates a name.

NOT NULL

Values of this domain are prevented from being null (but see notes below).

NULL

Values of this domain are allowed to be null. This is the default.

This clause is only intended for compatibility with nonstandard SQL databases. Its use is dis-
couraged in new applications.

CHECK (expression)

CHECK clauses specify integrity constraints or tests which values of the domain must satisfy. Each
constraint must be an expression producing a Boolean result. It should use the key word VALUE
to refer to the value being tested. Expressions evaluating to TRUE or UNKNOWN succeed. If
the expression produces a FALSE result, an error is reported and the value is not allowed to be
converted to the domain type.

Currently, CHECK expressions cannot contain subqueries nor refer to variables other than
VALUE.

When a domain has multiple CHECK constraints, they will be tested in alphabetical order by name.
(PostgreSQL versions before 9.5 did not honor any particular firing order for CHECK constraints.)

Notes

Domain constraints, particularly NOT NULL, are checked when converting a value to the domain type.
It is possible for a column that is nominally of the domain type to read as null despite there being
such a constraint. For example, this can happen in an outer-join query, if the domain column is on the
nullable side of the outer join. A more subtle example is

```
INSERT INTO tab (domcol) VALUES ((SELECT domcol FROM tab WHERE false));
```

The empty scalar sub-SELECT will produce a null value that is considered to be of the domain type,
so no further constraint checking is applied to it, and the insertion will succeed.

It is very difficult to avoid such problems, because of SQL's general assumption that a null value is a
valid value of every data type. Best practice therefore is to design a domain’s constraints so that a null
value is allowed, and then to apply column NOT NULL constraints to columns of the domain type as
needed, rather than directly to the domain type.

PostgreSQL assumes that CHECK constraints’ conditions are immutable, that is, they will always give
the same result for the same input value. This assumption is what justifies examining CHECK con-
straints only when a value is first converted to be of a domain type, and not at other times. (This is
essentially the same as the treatment of table CHECK constraints, as described in Section 5.3.1.)

An example of a common way to break this assumption is to reference a user-defined function in
a CHECK expression, and then change the behavior of that function. PostgreSQL does not disallow
CREATE DOMAIN

that, but it will not notice if there are stored values of the domain type that now violate the CHECK constraint. That would cause a subsequent database dump and reload to fail. The recommended way to handle such a change is to drop the constraint (using ALTER DOMAIN), adjust the function definition, and re-add the constraint, thereby rechecking it against stored data.

Examples

This example creates the us_postal_code data type and then uses the type in a table definition. A regular expression test is used to verify that the value looks like a valid US postal code:

CREATE DOMAIN us_postal_code AS TEXT
CHECK(
    VALUE ~ '^[0-9]{5}$'
    OR VALUE ~ '^[0-9]{5}-[0-9]{4}$'
);

CREATE TABLE us_snail_addy (  
    address_id SERIAL PRIMARY KEY,  
    street1 TEXT NOT NULL,  
    street2 TEXT,  
    street3 TEXT,  
    city TEXT NOT NULL,  
    postal us_postal_code NOT NULL  
);

Compatibility

The command CREATE DOMAIN conforms to the SQL standard.

See Also

ALTER DOMAIN, DROP DOMAIN
CREATE EVENT TRIGGER

Name
CREATE EVENT TRIGGER — define a new event trigger

Synopsis
CREATE EVENT TRIGGER name
  ON event
  [ WHEN filter_variable IN (filter_value [, ... ]) [ AND ... ] ]
  EXECUTE PROCEDURE function_name()

Description
CREATE EVENT TRIGGER creates a new event trigger. Whenever the designated event occurs and the WHEN condition associated with the trigger, if any, is satisfied, the trigger function will be executed. For a general introduction to event triggers, see Chapter 38. The user who creates an event trigger becomes its owner.

Parameters

name
The name to give the new trigger. This name must be unique within the database.

event
The name of the event that triggers a call to the given function. See Section 38.1 for more information on event names.

filter_variable
The name of a variable used to filter events. This makes it possible to restrict the firing of the trigger to a subset of the cases in which it is supported. Currently the only supported filter_variable is TAG.

filter_value
A list of values for the associated filter_variable for which the trigger should fire. For TAG, this means a list of command tags (e.g. ’DROP FUNCTION’).

function_name
A user-supplied function that is declared as taking no argument and returning type event_trigger.

Notes
Only superusers can create event triggers.
Event triggers are disabled in single-user mode (see postgres). If an erroneous event trigger disables the database so much that you can’t even drop the trigger, restart in single-user mode and you’ll be able to do that.

Examples

Forbid the execution of any DDL command:

```sql
CREATE OR REPLACE FUNCTION abort_any_command()
  RETURNS event_trigger
  LANGUAGE plpgsql
  AS $$
BEGIN
  RAISE EXCEPTION 'command % is disabled', tg_tag;
END;
$$;

CREATE EVENT TRIGGER abort_ddl ON ddl_command_start
  EXECUTE PROCEDURE abort_any_command();
```

Compatibility

There is no `CREATE EVENT TRIGGER` statement in the SQL standard.

See Also

`ALTER EVENT TRIGGER`, `DROP EVENT TRIGGER`, `CREATE FUNCTION`
CREATE EXTENSION

Name
CREATE EXTENSION — install an extension

Synopsis
CREATE EXTENSION [ IF NOT EXISTS ] extension_name
[ WITH ] [ SCHEMA schema_name ]
[ VERSION version ]
[ FROM old_version ]
[ CASCADE ]

Description
CREATE EXTENSION loads a new extension into the current database. There must not be an extension of the same name already loaded.

Loading an extension essentially amounts to running the extension’s script file. The script will typically create new SQL objects such as functions, data types, operators and index support methods. CREATE EXTENSION additionally records the identities of all the created objects, so that they can be dropped again if DROP EXTENSION is issued.

Loading an extension requires the same privileges that would be required to create its component objects. For most extensions this means superuser or database owner privileges are needed. The user who runs CREATE EXTENSION becomes the owner of the extension for purposes of later privilege checks, as well as the owner of any objects created by the extension’s script.

Parameters

IF NOT EXISTS

Do not throw an error if an extension with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing extension is anything like the one that would have been created from the currently-available script file.

extension_name

The name of the extension to be installed. PostgreSQL will create the extension using details from the file SHAREDIR/extension/extension_name.control.

schema_name

The name of the schema in which to install the extension’s objects, given that the extension allows its contents to be relocated. The named schema must already exist. If not specified, and the extension’s control file does not specify a schema either, the current default object creation schema is used.

If the extension specifies a schema parameter in its control file, then that schema cannot be overridden with a SCHEMA clause. Normally, an error will be raised if a SCHEMA clause is given and it conflicts with the extension’s schema parameter. However, if the CASCADE clause is also
given, then `schema_name` is ignored when it conflicts. The given `schema_name` will be used for installation of any needed extensions that do not specify `schema` in their control files.

Remember that the extension itself is not considered to be within any schema: extensions have unqualified names that must be unique database-wide. But objects belonging to the extension can be within schemas.

`version`

The version of the extension to install. This can be written as either an identifier or a string literal. The default version is whatever is specified in the extension’s control file.

`old_version`

`FROM old_version` must be specified when, and only when, you are attempting to install an extension that replaces an “old style” module that is just a collection of objects not packaged into an extension. This option causes `CREATE EXTENSION` to run an alternative installation script that absorbs the existing objects into the extension, instead of creating new objects. Be careful that `SCHEMA` specifies the schema containing these pre-existing objects.

The value to use for `old_version` is determined by the extension’s author, and might vary if there is more than one version of the old-style module that can be upgraded into an extension. For the standard additional modules supplied with pre-9.1 PostgreSQL, use `unpackaged` for `old_version` when updating a module to extension style.

`CASCADE`

Automatically install any extensions that this extension depends on that are not already installed. Their dependencies are likewise automatically installed, recursively. The `SCHEMA` clause, if given, applies to all extensions that get installed this way. Other options of the statement are not applied to automatically-installed extensions; in particular, their default versions are always selected.

**Notes**

Before you can use `CREATE EXTENSION` to load an extension into a database, the extension’s supporting files must be installed. Information about installing the extensions supplied with PostgreSQL can be found in Additional Supplied Modules.

The extensions currently available for loading can be identified from the `pg_available_extensions` or `pg_available_extension_versions` system views.

For information about writing new extensions, see Section 36.15.

**Examples**

Install the hstore extension into the current database:

```
CREATE EXTENSION hstore;
```

Update a pre-9.1 installation of hstore into extension style:

```
CREATE EXTENSION hstore SCHEMA public FROM unpackaged;
```

Be careful to specify the schema in which you installed the existing hstore objects.
Compatibility
CREATE EXTENSION is a PostgreSQL extension.

See Also
ALTER EXTENSION, DROP EXTENSION
CREATE FOREIGN DATA WRAPPER

Name

CREATE FOREIGN DATA WRAPPER — define a new foreign-data wrapper

Synopsis

CREATE FOREIGN DATA WRAPPER name
  [ HANDLER handler_function | NO HANDLER ]
  [ VALIDATOR validator_function | NO VALIDATOR ]
  [ OPTIONS ( option 'value' [, ... ] ) ]

Description

CREATE FOREIGN DATA WRAPPER creates a new foreign-data wrapper. The user who defines a foreign-data wrapper becomes its owner.

The foreign-data wrapper name must be unique within the database.

Only superusers can create foreign-data wrappers.

Parameters

name

The name of the foreign-data wrapper to be created.

HANDLER handler_function

handler_function is the name of a previously registered function that will be called to retrieve the execution functions for foreign tables. The handler function must take no arguments, and its return type must be fdw_handler.

It is possible to create a foreign-data wrapper with no handler function, but foreign tables using such a wrapper can only be declared, not accessed.

VALIDATOR validator_function

validator_function is the name of a previously registered function that will be called to check the generic options given to the foreign-data wrapper, as well as options for foreign servers, user mappings and foreign tables using the foreign-data wrapper. If no validator function or NO VALIDATOR is specified, then options will not be checked at creation time. (Foreign-data wrappers will possibly ignore or reject invalid option specifications at run time, depending on the implementation.) The validator function must take two arguments: one of type text[], which will contain the array of options as stored in the system catalogs, and one of type oid, which will be the OID of the system catalog containing the options. The return type is ignored; the function should report invalid options using the ereport(ERROR) function.

OPTIONS ( option 'value' [, ... ] )

This clause specifies options for the new foreign-data wrapper. The allowed option names and values are specific to each foreign data wrapper and are validated using the foreign-data wrapper’s validator function. Option names must be unique.
CREATE FOREIGN DATA WRAPPER

Notes

PostgreSQL’s foreign-data functionality is still under active development. Optimization of queries is primitive (and mostly left to the wrapper, too). Thus, there is considerable room for future performance improvements.

Examples

Create a useless foreign-data wrapper `dummy`:

```sql
CREATE FOREIGN DATA WRAPPER dummy;
```

Create a foreign-data wrapper `file` with handler function `file_fdw_handler`:

```sql
CREATE FOREIGN DATA WRAPPER file HANDLER file_fdw_handler;
```

Create a foreign-data wrapper `mywrapper` with some options:

```sql
CREATE FOREIGN DATA WRAPER mywrapper
    OPTIONS (debug 'true');
```

Compatibility

CREATE FOREIGN DATA WRAPPER conforms to ISO/IEC 9075-9 (SQL/MED), with the exception that the HANDLER and VALIDATOR clauses are extensions and the standard clauses LIBRARY and LANGUAGE are not implemented in PostgreSQL.

Note, however, that the SQL/MED functionality as a whole is not yet conforming.

See Also

ALTER FOREIGN DATA WRAPPER, DROP FOREIGN DATA WRAPPER, CREATE SERVER, CREATE USER MAPPING, CREATE FOREIGN TABLE
CREATE FOREIGN TABLE

Name

CREATE FOREIGN TABLE — define a new foreign table

Synopsis

CREATE FOREIGN TABLE [ IF NOT EXISTS ] table_name ( [ column_name data_type [ OPTIONS ( option 'value' [, ... ] ) ] [ COLLATE collation ] [ column_constraint ] [ ... ] ] [ INHERITS ( parent_table [, ... ] ) ] SERVER server_name [ OPTIONS ( option 'value' [, ... ] ) ]

where column_constraint is:

[ CONSTRAINT constraint_name ]
{ NOT NULL | NULL | CHECK ( expression ) [ NO INHERIT ] | DEFAULT default_expr }

and table_constraint is:

[ CONSTRAINT constraint_name ]
CHECK ( expression ) [ NO INHERIT ]

Description

CREATE FOREIGN TABLE creates a new foreign table in the current database. The table will be owned by the user issuing the command.

If a schema name is given (for example, CREATE FOREIGN TABLE myschema.mytable ...) then the table is created in the specified schema. Otherwise it is created in the current schema. The name of the foreign table must be distinct from the name of any other foreign table, table, sequence, index, view, or materialized view in the same schema.

CREATE FOREIGN TABLE also automatically creates a data type that represents the composite type corresponding to one row of the foreign table. Therefore, foreign tables cannot have the same name as any existing data type in the same schema.

To be able to create a foreign table, you must have USAGE privilege on the foreign server, as well as USAGE privilege on all column types used in the table.
CREATE FOREIGN TABLE

Parameters

IF NOT EXISTS
Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing relation is anything like the one that would have been created.

table_name
The name (optionally schema-qualified) of the table to be created.

column_name
The name of a column to be created in the new table.

data_type
The data type of the column. This can include array specifiers. For more information on the data types supported by PostgreSQL, refer to Chapter 8.

COLLATE collation
The COLLATE clause assigns a collation to the column (which must be of a collatable data type). If not specified, the column data type’s default collation is used.

INHERITS ( parent_table [, ... ] )
The optional INHERITS clause specifies a list of tables from which the new foreign table automatically inherits all columns. Parent tables can be plain tables or foreign tables. See the similar form of CREATE TABLE for more details.

CONSTRAINT constraint_name
An optional name for a column or table constraint. If the constraint is violated, the constraint name is present in error messages, so constraint names like col must be positive can be used to communicate helpful constraint information to client applications. (Double-quotes are needed to specify constraint names that contain spaces.) If a constraint name is not specified, the system generates a name.

NOT NULL
The column is not allowed to contain null values.

NULL
The column is allowed to contain null values. This is the default.

This clause is only provided for compatibility with non-standard SQL databases. Its use is discouraged in new applications.

CHECK ( expression ) [ NO INHERIT ]
The CHECK clause specifies an expression producing a Boolean result which each row in the foreign table is expected to satisfy; that is, the expression should produce TRUE or UNKNOWN, never FALSE, for all rows in the foreign table. A check constraint specified as a column constraint should reference that column’s value only, while an expression appearing in a table constraint can reference multiple columns.

Currently, CHECK expressions cannot contain subqueries nor refer to variables other than columns of the current row. The system column tableoid may be referenced, but not any other system column.

A constraint marked with NO INHERIT will not propagate to child tables.
DEFAULT default_expr

The DEFAULT clause assigns a default data value for the column whose column definition it appears within. The value is any variable-free expression (subqueries and cross-references to other columns in the current table are not allowed). The data type of the default expression must match the data type of the column.

The default expression will be used in any insert operation that does not specify a value for the column. If there is no default for a column, then the default is null.

server_name

The name of an existing foreign server to use for the foreign table. For details on defining a server, see CREATE SERVER.

OPTIONS ( option 'value' [, ...] )

Options to be associated with the new foreign table or one of its columns. The allowed option names and values are specific to each foreign data wrapper and are validated using the foreign-data wrapper’s validator function. Duplicate option names are not allowed (although it’s OK for a table option and a column option to have the same name).

Notes

Constraints on foreign tables (such as CHECK or NOT NULL clauses) are not enforced by the core PostgreSQL system, and most foreign data wrappers do not attempt to enforce them either; that is, the constraint is simply assumed to hold true. There would be little point in such enforcement since it would only apply to rows inserted or updated via the foreign table, and not to rows modified by other means, such as directly on the remote server. Instead, a constraint attached to a foreign table should represent a constraint that is being enforced by the remote server.

Some special-purpose foreign data wrappers might be the only access mechanism for the data they access, and in that case it might be appropriate for the foreign data wrapper itself to perform constraint enforcement. But you should not assume that a wrapper does that unless its documentation says so.

Although PostgreSQL does not attempt to enforce constraints on foreign tables, it does assume that they are correct for purposes of query optimization. If there are rows visible in the foreign table that do not satisfy a declared constraint, queries on the table might produce incorrect answers. It is the user’s responsibility to ensure that the constraint definition matches reality.

Examples

Create foreign table films, which will be accessed through the server film_server:

```sql
CREATE FOREIGN TABLE films (   code char(5) NOT NULL,   title varchar(40) NOT NULL,   did integer NOT NULL,   date_prod date,   kind varchar(10),   len interval hour to minute )   SERVER film_server;
```
Compatibility

The `CREATE FOREIGN TABLE` command largely conforms to the SQL standard; however, much as with `CREATE TABLE`, `NULL` constraints and zero-column foreign tables are permitted. The ability to specify column default values is also a PostgreSQL extension. Table inheritance, in the form defined by PostgreSQL, is nonstandard.

See Also

`ALTER FOREIGN TABLE`, `DROP FOREIGN TABLE`, `CREATE TABLE`, `CREATE SERVER`, `IMPORT FOREIGN SCHEMA`
CREATE FUNCTION

Name

CREATE FUNCTION — define a new function

Synopsis

CREATE [ OR REPLACE ] FUNCTION
  name ( [ [ argmode ] [ argname ] argtype [ { DEFAULT | = } default_expr ] [, ... ] ] )
  [ RETURNS rettype
    | RETURNS TABLE ( column_name column_type [, ... ] ) ]
  { LANGUAGE lang_name
    | TRANSFORM { FOR TYPE type_name } [, ... ]
    | WINDOW
    | IMMUTABLE | STABLE | VOLATILE | [ NOT ] LEAKPROOF
    | CALLED ON NULL INPUT | RETURNS NULL ON NULL INPUT | STRICT
    | [ EXTERNAL | SECURITY INVOKER | [ EXTERNAL ] SECURITY DEFINER
    | PARALLEL { UNSAFE | RESTRICTED | SAFE }
    | COST execution_cost
    | ROWS result_rows
    | SET configuration_parameter { TO value | = value | FROM CURRENT }
    | AS 'definition'
    | AS 'obj_file', 'link_symbol'
  } ...
  [ WITH ( attribute [, ... ] ) ]

Description

CREATE FUNCTION defines a new function. CREATE OR REPLACE FUNCTION will either create a new function, or replace an existing definition. To be able to define a function, the user must have the USAGE privilege on the language.

If a schema name is included, then the function is created in the specified schema. Otherwise it is created in the current schema. The name of the new function must not match any existing function with the same input argument types in the same schema. However, functions of different argument types can share a name (this is called overloading).

To replace the current definition of an existing function, use CREATE OR REPLACE FUNCTION. It is not possible to change the name or argument types of a function this way (if you tried, you would actually be creating a new, distinct function). Also, CREATE OR REPLACE FUNCTION will not let you change the return type of an existing function. To do that, you must drop and recreate the function. (When using OUT parameters, that means you cannot change the types of any OUT parameters except by dropping the function.)

When CREATE OR REPLACE FUNCTION is used to replace an existing function, the ownership and permissions of the function do not change. All other function properties are assigned the values specified or implied in the command. You must own the function to replace it (this includes being a member of the owning role).

If you drop and then recreate a function, the new function is not the same entity as the old; you will have to drop existing rules, views, triggers, etc. that refer to the old function. Use CREATE OR
CREATE FUNCTION

REPLACE FUNCTION to change a function definition without breaking objects that refer to the function. Also, ALTER FUNCTION can be used to change most of the auxiliary properties of an existing function.

The user that creates the function becomes the owner of the function.

To be able to create a function, you must have USAGE privilege on the argument types and the return type.

Parameters

name

The name (optionally schema-qualified) of the function to create.

argmode

The mode of an argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Only OUT arguments can follow a VARIADIC one. Also, OUT and INOUT arguments cannot be used together with the RETURNS TABLE notation.

argname

The name of an argument. Some languages (including SQL and PL/pgSQL) let you use the name in the function body. For other languages the name of an input argument is just extra documentation, so far as the function itself is concerned; but you can use input argument names when calling a function to improve readability (see Section 4.3). In any case, the name of an output argument is significant, because it defines the column name in the result row type. (If you omit the name for an output argument, the system will choose a default column name.)

argtype

The data type(s) of the function’s arguments (optionally schema-qualified), if any. The argument types can be base, composite, or domain types, or can reference the type of a table column.

Depending on the implementation language it might also be allowed to specify “pseudotypes” such as cstring. Pseudotypes indicate that the actual argument type is either incompletely specified, or outside the set of ordinary SQL data types.

The type of a column is referenced by writing table_name.column_name%TYPE. Using this feature can sometimes help make a function independent of changes to the definition of a table.

default_expr

An expression to be used as default value if the parameter is not specified. The expression has to be coercible to the argument type of the parameter. Only input (including INOUT) parameters can have a default value. All input parameters following a parameter with a default value must have default values as well.

rertype

The return data type (optionally schema-qualified). The return type can be a base, composite, or domain type, or can reference the type of a table column. Depending on the implementation language it might also be allowed to specify “pseudotypes” such as cstring. If the function is not supposed to return a value, specify void as the return type.

When there are OUT or INOUT parameters, the RETURNS clause can be omitted. If present, it must agree with the result type implied by the output parameters: RECORD if there are multiple output parameters, or the same type as the single output parameter.
The SETOF modifier indicates that the function will return a set of items, rather than a single item.

The type of a column is referenced by writing `table_name.column_name%TYPE`.

`column_name`  
The name of an output column in the `RETURNS TABLE` syntax. This is effectively another way of declaring a named `OUT` parameter, except that `RETURNS TABLE` also implies `RETURNS SETOF`.

`column_type`  
The data type of an output column in the `RETURNS TABLE` syntax.

`lang_name`  
The name of the language that the function is implemented in. It can be `sql`, `c`, `internal`, or the name of a user-defined procedural language, e.g. `plpgsql`. Enclosing the name in single quotes is deprecated and requires matching case.

`TRANSFORM { FOR TYPE type_name | [, ... ] }`  
Lists which transforms a call to the function should apply. Transforms convert between SQL types and language-specific data types; see CREATE TRANSFORM. Procedural language implementations usually have hardcoded knowledge of the built-in types, so those don’t need to be listed here. If a procedural language implementation does not know how to handle a type and no transform is supplied, it will fall back to a default behavior for converting data types, but this depends on the implementation.

`WINDOW`  
`WINDOW` indicates that the function is a window function rather than a plain function. This is currently only useful for functions written in C. The `WINDOW` attribute cannot be changed when replacing an existing function definition.

`IMMUTABLE`  
`STABLE`  
`VOLATILE`  
These attributes inform the query optimizer about the behavior of the function. At most one choice can be specified. If none of these appear, `VOLATILE` is the default assumption.

`IMMUTABLE` indicates that the function cannot modify the database and always returns the same result when given the same argument values; that is, it does not do database lookups or otherwise use information not directly present in its argument list. If this option is given, any call of the function with all-constant arguments can be immediately replaced with the function value.

`STABLE` indicates that the function cannot modify the database, and that within a single table scan it will consistently return the same result for the same argument values, but that its result could change across SQL statements. This is the appropriate selection for functions whose results depend on database lookups, parameter variables (such as the current time zone), etc. (It is inappropriate for AFTER triggers that wish to query rows modified by the current command.) Also note that the `current_timestamp` family of functions qualify as stable, since their values do not change within a transaction.

`VOLATILE` indicates that the function value can change even within a single table scan, so no optimizations can be made. Relatively few database functions are volatile in this sense; some examples are `random()`, `currval()`, `timeofday()`. But note that any function that has side-effects must be classified volatile, even if its result is quite predictable, to prevent calls from being optimized away; an example is `setval()`.

For additional details see Section 36.6.
CREATE FUNCTION

LEAKPROOF

**LEAKPROOF** indicates that the function has no side effects. It reveals no information about its arguments other than by its return value. For example, a function which throws an error message for some argument values but not others, or which includes the argument values in any error message, is not leakproof. This affects how the system executes queries against views created with the **security_barrier** option or tables with row level security enabled. The system will enforce conditions from security policies and security barrier views before any user-supplied conditions from the query itself that contain non-leakproof functions, in order to prevent the inadvertent exposure of data. Functions and operators marked as leakproof are assumed to be trustworthy, and may be executed before conditions from security policies and security barrier views. In addition, functions which do not take arguments or which are not passed any arguments from the security barrier view or table do not have to be marked as leakproof to be executed before security conditions. See CREATE VIEW and Section 39.5. This option can only be set by the superuser.

**CALLED ON NULL INPUT**

**RETURNS NULL ON NULL INPUT**

**STRICT**

**CALLED ON NULL INPUT** (the default) indicates that the function will be called normally when some of its arguments are null. It is then the function author’s responsibility to check for null values if necessary and respond appropriately.

**RETURNS NULL ON NULL INPUT or STRICT** indicates that the function always returns null whenever any of its arguments are null. If this parameter is specified, the function is not executed when there are null arguments; instead a null result is assumed automatically.

**[EXTERNAL] SECURITY INVOKER**

**[EXTERNAL] SECURITY DEFINER**

**SECURITY INVOKER** indicates that the function is to be executed with the privileges of the user that calls it. That is the default. **SECURITY DEFINER** specifies that the function is to be executed with the privileges of the user that created it.

The key word **EXTERNAL** is allowed for SQL conformance, but it is optional since, unlike in SQL, this feature applies to all functions not only external ones.

**PARALLEL**

**PARALLEL UNSAFE** indicates that the function can’t be executed in parallel mode and the presence of such a function in an SQL statement forces a serial execution plan. This is the default. **PARALLEL RESTRICTED** indicates that the function can be executed in parallel mode, but the execution is restricted to parallel group leader. **PARALLEL SAFE** indicates that the function is safe to run in parallel mode without restriction.

Functions should be labeled parallel unsafe if they modify any database state, or if they make changes to the transaction such as using sub-transactions, or if they access sequences or attempt to make persistent changes to settings (e.g. **setval**). They should be labeled as parallel restricted if they access temporary tables, client connection state, cursors, prepared statements, or miscellaneous backend-local state which the system cannot synchronize in parallel mode (e.g. **setseed** cannot be executed other than by the group leader because a change made by another process would not be reflected in the leader). In general, if a function is labeled as being safe when it is restricted or unsafe, or if it is labeled as being restricted when it is in fact unsafe, it may throw errors or produce wrong answers when used in a parallel query. C-language functions could in theory exhibit totally undefined behavior if mislabeled, since there is no way for the system to protect itself against arbitrary C code, but in most likely cases the result will be no worse than for any other function. If in doubt, functions should be labeled as **UNSAFE**, which is the default.
CREATE FUNCTION

execution_cost
A positive number giving the estimated execution cost for the function, in units of cpu_operator_cost. If the function returns a set, this is the cost per returned row. If the cost is not specified, 1 unit is assumed for C-language and internal functions, and 100 units for functions in all other languages. Larger values cause the planner to try to avoid evaluating the function more often than necessary.

result_rows
A positive number giving the estimated number of rows that the planner should expect the function to return. This is only allowed when the function is declared to return a set. The default assumption is 1000 rows.

configuration_parameter
value
The SET clause causes the specified configuration parameter to be set to the specified value when the function is entered, and then restored to its prior value when the function exits. SET FROM CURRENT saves the value of the parameter that is current when CREATE FUNCTION is executed as the value to be applied when the function is entered.

If a SET clause is attached to a function, then the effects of a SET LOCAL command executed inside the function for the same variable are restricted to the function: the configuration parameter’s prior value is still restored at function exit. However, an ordinary SET command (without LOCAL) overrides the SET clause, much as it would do for a previous SET LOCAL command: the effects of such a command will persist after function exit, unless the current transaction is rolled back.

See SET and Chapter 19 for more information about allowed parameter names and values.

definition
A string constant defining the function; the meaning depends on the language. It can be an internal function name, the path to an object file, an SQL command, or text in a procedural language.

It is often helpful to use dollar quoting (see Section 4.1.2.4) to write the function definition string, rather than the normal single quote syntax. Without dollar quoting, any single quotes or backslashes in the function definition must be escaped by doubling them.

obj_file, link_symbol
This form of the AS clause is used for dynamically loadable C language functions when the function name in the C language source code is not the same as the name of the SQL function. The string obj_file is the name of the file containing the dynamically loadable object, and link_symbol is the function’s link symbol, that is, the name of the function in the C language source code. If the link symbol is omitted, it is assumed to be the same as the name of the SQL function being defined. The C names of all functions must be different, so you must give overloaded C functions different C names (for example, use the argument types as part of the C names).

When repeated CREATE FUNCTION calls refer to the same object file, the file is only loaded once per session. To unload and reload the file (perhaps during development), start a new session.

attribute
The historical way to specify optional pieces of information about the function. The following
CREATE FUNCTION

attributes can appear here:

isStrict

Equivalent to STRICT or RETURNS NULL ON NULL INPUT.

isCachable

isCachable is an obsolete equivalent of IMMUTABLE; it’s still accepted for backwards-compatibility reasons.

Attribute names are not case-sensitive.

Refer to Section 36.3 for further information on writing functions.

Overloading

PostgreSQL allows function overloading; that is, the same name can be used for several different functions so long as they have distinct input argument types. Whether or not you use it, this capability entails security precautions when calling functions in databases where some users mistrust other users; see Section 10.3.

Two functions are considered the same if they have the same names and input argument types, ignoring any OUT parameters. Thus for example these declarations conflict:

CREATE FUNCTION foo(int) ...
CREATE FUNCTION foo(int, out text) ...

Functions that have different argument type lists will not be considered to conflict at creation time, but if defaults are provided they might conflict in use. For example, consider

CREATE FUNCTION foo(int) ...
CREATE FUNCTION foo(int, int default 42) ...

A call foo(10) will fail due to the ambiguity about which function should be called.

Notes

The full SQL type syntax is allowed for declaring a function’s arguments and return value. However, parenthesized type modifiers (e.g., the precision field for type numeric) are discarded by CREATE FUNCTION. Thus for example CREATE FUNCTION foo (varchar(10)) ... is exactly the same as CREATE FUNCTION foo (varchar) ....

When replacing an existing function with CREATE OR REPLACE FUNCTION, there are restrictions on changing parameter names. You cannot change the name already assigned to any input parameter (although you can add names to parameters that had none before). If there is more than one output parameter, you cannot change the names of the output parameters, because that would change the column names of the anonymous composite type that describes the function’s result. These restrictions are made to ensure that existing calls of the function do not stop working when it is replaced.

If a function is declared STRICT with a VARIADIC argument, the strictness check tests that the variadic array as a whole is non-null. The function will still be called if the array has null elements.
CREATE FUNCTION Examples

Here are some trivial examples to help you get started. For more information and examples, see Section 36.3.

CREATE FUNCTION add(integer, integer) RETURNS integer
  AS 'select $1 + $2;'
  LANGUAGE SQL
  IMMUTABLE
  RETURNS NULL ON NULL INPUT;

Increment an integer, making use of an argument name, in PL/pgSQL:

CREATE OR REPLACE FUNCTION increment(i integer) RETURNS integer AS $$
BEGIN
  RETURN i + 1;
END;
$$ LANGUAGE plpgsql;

Return a record containing multiple output parameters:

CREATE FUNCTION dup(in int, out f1 int, out f2 text)
  AS $$ SELECT $1, CAST($1 AS text) || ' is text' $$
  LANGUAGE SQL;

SELECT * FROM dup(42);

You can do the same thing more verbosely with an explicitly named composite type:

CREATE TYPE dup_result AS (f1 int, f2 text);

CREATE FUNCTION dup(int) RETURNS dup_result
  AS $$ SELECT $1, CAST($1 AS text) || ' is text' $$
  LANGUAGE SQL;

SELECT * FROM dup(42);

Another way to return multiple columns is to use a TABLE function:

CREATE FUNCTION dup(int) RETURNS TABLE(f1 int, f2 text)
  AS $$ SELECT $1, CAST($1 AS text) || ' is text' $$
  LANGUAGE SQL;

SELECT * FROM dup(42);

However, a TABLE function is different from the preceding examples, because it actually returns a set of records, not just one record.

Writing SECURITY DEFINER Functions Safely

Because a SECURITY DEFINER function is executed with the privileges of the user that created it, care is needed to ensure that the function cannot be misused. For security, search_path should be set to exclude any schemas writable by untrusted users. This prevents malicious users from creating
CREATE FUNCTION objects (e.g., tables, functions, and operators) that mask objects intended to be used by the function. Particularly important in this regard is the temporary-table schema, which is searched first by default, and is normally writable by anyone. A secure arrangement can be obtained by forcing the temporary schema to be searched last. To do this, write `pg_temp` as the last entry in `search_path`. This function illustrates safe usage:

```sql
CREATE FUNCTION check_password(uname TEXT, pass TEXT)
RETURNS BOOLEAN AS $$
DECLARE passed BOOLEAN;
BEGIN
    SELECT (pwd = $2) INTO passed
    FROM pwds
    WHERE username = $1;
    RETURN passed;
END;
$$ LANGUAGE plpgsql
SECURITY DEFINER
-- Set a secure search_path: trusted schema(s), then 'pg_temp'.
SET search_path = admin, pg_temp;
```

This function's intention is to access a table `admin.pwds`. But without the `SET` clause, or with a `SET` clause mentioning only `admin`, the function could be subverted by creating a temporary table named `pwds`.

Before PostgreSQL version 8.3, the `SET` clause was not available, and so older functions may contain rather complicated logic to save, set, and restore `search_path`. The `SET` clause is far easier to use for this purpose.

Another point to keep in mind is that by default, execute privilege is granted to `PUBLIC` for newly created functions (see GRANT for more information). Frequently you will wish to restrict use of a security definer function to only some users. To do that, you must revoke the default `PUBLIC` privileges and then grant execute privilege selectively. To avoid having a window where the new function is accessible to all, create it and set the privileges within a single transaction. For example:

```sql
BEGIN;
CREATE FUNCTION check_password(uname TEXT, pass TEXT) ... SECURITY DEFINER;
REVOKE ALL ON FUNCTION check_password(uname TEXT, pass TEXT) FROM PUBLIC;
GRANT EXECUTE ON FUNCTION check_password(uname TEXT, pass TEXT) TO admins;
COMMIT;
```

**Compatibility**

A `CREATE FUNCTION` command is defined in SQL:1999 and later. The PostgreSQL version is similar but not fully compatible. The attributes are not portable, neither are the different available languages.

For compatibility with some other database systems, `argmode` can be written either before or after `argname`. But only the first way is standard-compliant.

For parameter defaults, the SQL standard specifies only the syntax with the `DEFAULT` key word. The syntax with `=` is used in T-SQL and Firebird.
See Also

ALTER FUNCTION, DROP FUNCTION, GRANT, LOAD, REVOKE, createlang
CREATE GROUP

Name
CREATE GROUP — define a new database role

Synopsis
CREATE GROUP name [ [ WITH ] option [ ... ] ]

where option can be:

- SUPERUSER | NOSUPERUSER
- CREATEDB | NOCREATEDB
- CREATEROLE | NOCREATEROLE
- INHERIT | NOINHERIT
- LOGIN | NOLOGIN
- [ ENCRIPTED | UNENCRYPTED ] PASSWORD ‘password’
- VALID UNTIL ‘timestamp’
- IN ROLE role_name [, ...]
- IN GROUP role_name [, ...]
- ROLE role_name [, ...]
- ADMIN role_name [, ...]
- USER role_name [, ...]
- SYSID uid

Description
CREATE GROUP is now an alias for CREATE ROLE.

Compatibility
There is no CREATE GROUP statement in the SQL standard.

See Also
CREATE ROLE
CREATE INDEX

Name

CREATE INDEX — define a new index

Synopsis

CREATE [ UNIQUE ] INDEX [ CONCURRENTLY ] [ [ IF NOT EXISTS ] name ] ON table_name [ USING method ]
( { column_name | ( expression ) } [ COLLATE collation ] [ opclass ] [ ASC | DESC ] [ NULLS { FIRST | LAST } ] [ WITH ( storage_parameter = value [, ... ] ) ]
[ TABLESPACE tablespace_name ]
[ WHERE predicate ]

Description

CREATE INDEX constructs an index on the specified column(s) of the specified relation, which can be a table or a materialized view. Indexes are primarily used to enhance database performance (though inappropriate use can result in slower performance).

The key field(s) for the index are specified as column names, or alternatively as expressions written in parentheses. Multiple fields can be specified if the index method supports multicolumn indexes.

An index field can be an expression computed from the values of one or more columns of the table row. This feature can be used to obtain fast access to data based on some transformation of the basic data. For example, an index computed on \( \text{upper}(\text{col}) \) would allow the clause \( \text{WHERE} \ \text{upper}(\text{col}) = \text{’JIM’} \) to use an index.

PostgreSQL provides the index methods B-tree, hash, GiST, SP-GiST, GIN, and BRIN. Users can also define their own index methods, but that is fairly complicated.

When the \text{WHERE} clause is present, a \text{partial index} is created. A partial index is an index that contains entries for only a portion of a table, usually a portion that is more useful for indexing than the rest of the table. For example, if you have a table that contains both billed and unbilled orders where the unbilled orders take up a small fraction of the total table and yet that is an often used section, you can improve performance by creating an index on just that portion. Another possible application is to use \text{WHERE} with \texttt{UNIQUE} to enforce uniqueness over a subset of a table. See Section 11.8 for more discussion.

The expression used in the \text{WHERE} clause can refer only to columns of the underlying table, but it can use all columns, not just the ones being indexed. Presently, subqueries and aggregate expressions are also forbidden in \text{WHERE}. The same restrictions apply to index fields that are expressions.

All functions and operators used in an index definition must be “immutable”, that is, their results must depend only on their arguments and never on any outside influence (such as the contents of another table or the current time). This restriction ensures that the behavior of the index is well-defined. To use a user-defined function in an index expression or \text{WHERE} clause, remember to mark the function immutable when you create it.
Parameters

UNIQUE

Causes the system to check for duplicate values in the table when the index is created (if data already exist) and each time data is added. Attempts to insert or update data which would result in duplicate entries will generate an error.

CONCURRENTLY

When this option is used, PostgreSQL will build the index without taking any locks that prevent concurrent inserts, updates, or deletes on the table; whereas a standard index build locks out writes (but not reads) on the table until it’s done. There are several caveats to be aware of when using this option — see Building Indexes Concurrently.

For temporary tables, CREATE INDEX is always non-concurrent, as no other session can access them, and non-concurrent index creation is cheaper.

IF NOT EXISTS

Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing index is anything like the one that would have been created. Index name is required when IF NOT EXISTS is specified.

name

The name of the index to be created. No schema name can be included here; the index is always created in the same schema as its parent table. If the name is omitted, PostgreSQL chooses a suitable name based on the parent table’s name and the indexed column name(s).

table_name

The name (possibly schema-qualified) of the table to be indexed.

method

The name of the index method to be used. Choices are btree, hash, gist, spgist, gin, and brin. The default method is btree.

column_name

The name of a column of the table.

expression

An expression based on one or more columns of the table. The expression usually must be written with surrounding parentheses, as shown in the syntax. However, the parentheses can be omitted if the expression has the form of a function call.

collation

The name of the collation to use for the index. By default, the index uses the collation declared for the column to be indexed or the result collation of the expression to be indexed. Indexes with non-default collations can be useful for queries that involve expressions using non-default collations.

opclass

The name of an operator class. See below for details.

ASC

Specifies ascending sort order (which is the default).
CREATE INDEX

desc
  Specifies descending sort order.

collate NULLS FIRST
  Specifies that nulls sort before non-nulls. This is the default when desc is specified.

collate NULLS LAST
  Specifies that nulls sort after non-nulls. This is the default when desc is not specified.

storage_parameter
  The name of an index-method-specific storage parameter. See Index Storage Parameters for details.

tablespace_name
  The tablespace in which to create the index. If not specified, default_tablespace is consulted, or temp_tablespaces for indexes on temporary tables.

predicate
  The constraint expression for a partial index.

Index Storage Parameters

The optional WITH clause specifies storage parameters for the index. Each index method has its own set of allowed storage parameters. The B-tree, hash, GiST and SP-GiST index methods all accept this parameter:

fillfactor
  The fillfactor for an index is a percentage that determines how full the index method will try to pack index pages. For B-trees, leaf pages are filled to this percentage during initial index build, and also when extending the index at the right (adding new largest key values). If pages subsequently become completely full, they will be split, leading to gradual degradation in the index’s efficiency. B-trees use a default fillfactor of 90, but any integer value from 10 to 100 can be selected. If the table is static then fillfactor 100 is best to minimize the index’s physical size, but for heavily updated tables a smaller fillfactor is better to minimize the need for page splits. The other index methods use fillfactor in different but roughly analogous ways; the default fillfactor varies between methods.

GiST indexes additionally accept this parameter:

buffering
  Determines whether the buffering build technique described in Section 61.4.1 is used to build the index. With off it is disabled, with on it is enabled, and with auto it is initially disabled, but turned on-the-fly once the index size reaches effective_cache_size. The default is auto.

GIN indexes accept different parameters:

fastupdate
  This setting controls usage of the fast update technique described in Section 63.4.1. It is a Boolean parameter: on enables fast update, off disables it. (Alternative spellings of on and off are allowed as described in Section 19.1.) The default is on.

Note: Turning fastupdate off via ALTER INDEX prevents future insertions from going into the list of pending index entries, but does not in itself flush previous entries. You might want to
CREATE INDEX

VACUUM the table or call gin_clean_pending_list function afterward to ensure the pending list is emptied.

gin_pending_list_limit

Custom gin_pending_list_limit parameter. This value is specified in kilobytes.

BRIN indexes accept a different parameter:

pages_per_range

Defines the number of table blocks that make up one block range for each entry of a BRIN index (see Section 64.1 for more details). The default is 128.

Building Indexes Concurrently

Creating an index can interfere with regular operation of a database. Normally PostgreSQL locks the table to be indexed against writes and performs the entire index build with a single scan of the table. Other transactions can still read the table, but if they try to insert, update, or delete rows in the table they will block until the index build is finished. This could have a severe effect if the system is a live production database. Very large tables can take many hours to be indexed, and even for smaller tables, an index build can lock out writers for periods that are unacceptably long for a production system.

PostgreSQL supports building indexes without locking out writes. This method is invoked by specifying the CONCURRENTLY option of CREATE INDEX. When this option is used, PostgreSQL must perform two scans of the table, and in addition it must wait for all existing transactions that could potentially modify or use the index to terminate. Thus this method requires more total work than a standard index build and takes significantly longer to complete. However, since it allows normal operations to continue while the index is built, this method is useful for adding new indexes in a production environment. Of course, the extra CPU and I/O load imposed by the index creation might slow other operations.

In a concurrent index build, the index is actually entered into the system catalogs in one transaction, then two table scans occur in two more transactions. Before each table scan, the index build must wait for existing transactions that have modified the table to terminate. After the second scan, the index build must wait for any transactions that have a snapshot (see Chapter 13) predating the second scan to terminate. Then finally the index can be marked ready for use, and the CREATE INDEX command terminates. Even then, however, the index may not be immediately usable for queries: in the worst case, it cannot be used as long as transactions exist that predate the start of the index build.

If a problem arises while scanning the table, such as a deadlock or a uniqueness violation in a unique index, the CREATE INDEX command will fail but leave behind an “invalid” index. This index will be ignored for querying purposes because it might be incomplete; however it will still consume update overhead. The psql \d command will report such an index as INVALID:

postgres=# \d tab
  Table "public.tab"
  Column | Type    | Modifiers
----------|---------|-----------
col      | integer |-----------
Indexes:
  "idx" btree (col) INVALID
The recommended recovery method in such cases is to drop the index and try again to perform `CREATE INDEX CONCURRENTLY`. (Another possibility is to rebuild the index with `REINDEX`. However, since `REINDEX` does not support concurrent builds, this option is unlikely to seem attractive.)

Another caveat when building a unique index concurrently is that the uniqueness constraint is already being enforced against other transactions when the second table scan begins. This means that constraint violations could be reported in other queries prior to the index becoming available for use, or even in cases where the index build eventually fails. Also, if a failure does occur in the second scan, the “invalid” index continues to enforce its uniqueness constraint afterwards.

Concurrent builds of expression indexes and partial indexes are supported. Errors occurring in the evaluation of these expressions could cause behavior similar to that described above for unique constraint violations.

Regular index builds permit other regular index builds on the same table to occur in parallel, but only one concurrent index build can occur on a table at a time. In both cases, no other types of schema modification on the table are allowed meanwhile. Another difference is that a regular `CREATE INDEX` command can be performed within a transaction block, but `CREATE INDEX CONCURRENTLY` cannot.

**Notes**

See Chapter 11 for information about when indexes can be used, when they are not used, and in which particular situations they can be useful.

---

**Caution**

Hash index operations are not presently WAL-logged, so hash indexes might need to be rebuilt with `REINDEX` after a database crash if there were unwritten changes. Also, changes to hash indexes are not replicated over streaming or file-based replication after the initial base backup, so they give wrong answers to queries that subsequently use them. Hash indexes are also not properly restored during point-in-time recovery. For these reasons, hash index use is presently discouraged.

Currently, only the B-tree, GiST, GIN, and BRIN index methods support multicolumn indexes. Up to 32 fields can be specified by default. (This limit can be altered when building PostgreSQL.) Only B-tree currently supports unique indexes.

An operator class can be specified for each column of an index. The operator class identifies the operators to be used by the index for that column. For example, a B-tree index on four-byte integers would use the `int4_ops` class; this operator class includes comparison functions for four-byte integers. In practice the default operator class for the column’s data type is usually sufficient. The main point of having operator classes is that for some data types, there could be more than one meaningful ordering. For example, we might want to sort a complex-number data type either by absolute value or by real part. We could do this by defining two operator classes for the data type and then selecting the proper class when making an index. More information about operator classes is in Section 11.9 and in Section 36.14.

For index methods that support ordered scans (currently, only B-tree), the optional clauses `ASC`, `DESC`, `NULLS FIRST`, and/or `NULLS LAST` can be specified to modify the sort ordering of the index. Since an ordered index can be scanned either forward or backward, it is not normally useful to create a single-column `DESC` index — that sort ordering is already available with a regular index. The value of these options is that multicolumn indexes can be created that match the sort ordering requested by
CREATE INDEX

a mixed-ordering query, such as SELECT ... ORDER BY x ASC, y DESC. The NULLS options are useful if you need to support “nulls sort low” behavior, rather than the default “nulls sort high”, in queries that depend on indexes to avoid sorting steps.

For most index methods, the speed of creating an index is dependent on the setting of maintenance_work_mem. Larger values will reduce the time needed for index creation, so long as you don’t make it larger than the amount of memory really available, which would drive the machine into swapping.

Use DROP INDEX to remove an index.

Prior releases of PostgreSQL also had an R-tree index method. This method has been removed because it had no significant advantages over the GiST method. If USING rtree is specified, CREATE INDEX will interpret it as USING gist, to simplify conversion of old databases to GiST.

Examples

To create a B-tree index on the column title in the table films:

CREATE UNIQUE INDEX title_idx ON films (title);

To create an index on the expression lower(title), allowing efficient case-insensitive searches:

CREATE INDEX ON films ((lower(title)));

(In this example we have chosen to omit the index name, so the system will choose a name, typically films_lower_idx.)

To create an index with non-default collation:

CREATE INDEX title_idx_german ON films (title COLLATE "de_DE");

To create an index with non-default sort ordering of nulls:

CREATE INDEX title_idx_nulls_low ON films (title NULLS FIRST);

To create an index with non-default fill factor:

CREATE UNIQUE INDEX title_idx ON films (title) WITH (fillfactor = 70);

To create a GIN index with fast updates disabled:

CREATE INDEX gin_idx ON documents_table USING GIN (locations) WITH (fastupdate = off);

To create an index on the column code in the table films and have the index reside in the tablespace indexspace:

CREATE INDEX code_idx ON films (code) TABLESPACE indexspace;
CREATE INDEX

To create a GiST index on a point attribute so that we can efficiently use box operators on the result of the conversion function:

```
CREATE INDEX pointloc
  ON points USING gist (box(location,location));
SELECT * FROM points
  WHERE box(location,location) && '(0,0),(1,1)'::box;
```

To create an index without locking out writes to the table:

```
CREATE INDEX CONCURRENTLY sales_quantity_index ON sales_table (quantity);
```

**Compatibility**

CREATE INDEX is a PostgreSQL language extension. There are no provisions for indexes in the SQL standard.

**See Also**

ALTER INDEX, DROP INDEX
CREATE LANGUAGE

Name
CREATE LANGUAGE — define a new procedural language

Synopsis
CREATE [ OR REPLACE ] [ PROCEDURAL ] LANGUAGE name
CREATE [ OR REPLACE ] [ TRUSTED ] [ PROCEDURAL ] LANGUAGE name
    HANDLER call_handler [ INLINE inline_handler ] [ VALIDATOR valfunction ]

Description
CREATE LANGUAGE registers a new procedural language with a PostgreSQL database. Subsequently, functions and trigger procedures can be defined in this new language.

Note: As of PostgreSQL 9.1, most procedural languages have been made into “extensions”, and should therefore be installed with CREATE EXTENSION not CREATE LANGUAGE. Direct use of CREATE LANGUAGE should now be confined to extension installation scripts. If you have a “bare” language in your database, perhaps as a result of an upgrade, you can convert it to an extension using CREATE EXTENSION langname FROM unpackaged.

CREATE LANGUAGE effectively associates the language name with handler function(s) that are responsible for executing functions written in the language. Refer to Chapter 54 for more information about language handlers.

There are two forms of the CREATE LANGUAGE command. In the first form, the user supplies just the name of the desired language, and the PostgreSQL server consults the pg_pltemplate system catalog to determine the correct parameters. In the second form, the user supplies the language parameters along with the language name. The second form can be used to create a language that is not defined in pg_pltemplate, but this approach is considered obsolescent.

When the server finds an entry in the pg_pltemplate catalog for the given language name, it will use the catalog data even if the command includes language parameters. This behavior simplifies loading of old dump files, which are likely to contain out-of-date information about language support functions.

Ordinarily, the user must have the PostgreSQL superuser privilege to register a new language. However, the owner of a database can register a new language within that database if the language is listed in the pg_pltemplate catalog and is marked as allowed to be created by database owners (tmpldbacreate is true). The default is that trusted languages can be created by database owners, but this can be adjusted by superusers by modifying the contents of pg_pltemplate. The creator of a language becomes its owner and can later drop it, rename it, or assign it to a new owner.

CREATE OR REPLACE LANGUAGE will either create a new language, or replace an existing definition. If the language already exists, its parameters are updated according to the values specified or taken from pg_pltemplate, but the language’s ownership and permissions settings do not change, and any existing functions written in the language are assumed to still be valid. In addition to the normal privilege requirements for creating a language, the user must be superuser or owner of the existing language. The REPLACE case is mainly meant to be used to ensure that the language exists. If the
language has a pg_pltemplate entry then REPLACE will not actually change anything about an existing definition, except in the unusual case where the pg_pltemplate entry has been modified since the language was created.

Parameters

TRUSTED

TRUSTED specifies that the language does not grant access to data that the user would not otherwise have. If this key word is omitted when registering the language, only users with the PostgreSQL superuser privilege can use this language to create new functions.

PROCEDURAL

This is a noise word.

name

The name of the new procedural language. The name must be unique among the languages in the database.

For backward compatibility, the name can be enclosed by single quotes.

HANDLER call_handler

call_handler is the name of a previously registered function that will be called to execute the procedural language’s functions. The call handler for a procedural language must be written in a compiled language such as C with version 1 call convention and registered with PostgreSQL as a function taking no arguments and returning the language_handler type, a placeholder type that is simply used to identify the function as a call handler.

INLINE inline_handler

inline_handler is the name of a previously registered function that will be called to execute an anonymous code block (DO command) in this language. If no inline_handler function is specified, the language does not support anonymous code blocks. The handler function must take one argument of type internal, which will be the DO command’s internal representation, and it will typically return void. The return value of the handler is ignored.

VALIDATOR valfunction

valfunction is the name of a previously registered function that will be called when a new function in the language is created, to validate the new function. If no validator function is specified, then a new function will not be checked when it is created. The validator function must take one argument of type oid, which will be the OID of the to-be-created function, and will typically return void.

A validator function would typically inspect the function body for syntactical correctness, but it can also look at other properties of the function, for example if the language cannot handle certain argument types. To signal an error, the validator function should use the ereport() function. The return value of the function is ignored.

The TRUSTED option and the support function name(s) are ignored if the server has an entry for the specified language name in pg_pltemplate.
CREATE LANGUAGE

Notes

The createlang program is a simple wrapper around the CREATE LANGUAGE command. It eases installation of procedural languages from the shell command line.

Use DROP LANGUAGE, or better yet the droplang program, to drop procedural languages.

The system catalog pg_language (see Section 50.29) records information about the currently installed languages. Also, createlang has an option to list the installed languages.

To create functions in a procedural language, a user must have the USAGE privilege for the language. By default, USAGE is granted to PUBLIC (i.e., everyone) for trusted languages. This can be revoked if desired.

Procedural languages are local to individual databases. However, a language can be installed into the template1 database, which will cause it to be available automatically in all subsequently-created databases.

The call handler function, the inline handler function (if any), and the validator function (if any) must already exist if the server does not have an entry for the language in pg_pltemplate. But when there is an entry, the functions need not already exist; they will be automatically defined if not present in the database. (This might result in CREATE LANGUAGE failing, if the shared library that implements the language is not available in the installation.)

In PostgreSQL versions before 7.3, it was necessary to declare handler functions as returning the placeholder type opaque, rather than language_handler. To support loading of old dump files, CREATE LANGUAGE will accept a function declared as returning opaque, but it will issue a notice and change the function’s declared return type to language_handler.

Examples

The preferred way of creating any of the standard procedural languages is just:

```
CREATE LANGUAGE plperl;
```

For a language not known in the pg_pltemplate catalog, a sequence such as this is needed:

```
CREATE FUNCTION plsample_call_handler() RETURNS language_handler
    AS '$libdir/plsample'
    LANGUAGE C;
CREATE LANGUAGE plsample
    HANDLER plsample_call_handler;
```

Compatibility

CREATE LANGUAGE is a PostgreSQL extension.

See Also

ALTER LANGUAGE, CREATE FUNCTION, DROP LANGUAGE, GRANT, REVOKE, createlang, droplang
CREATE MATERIALIZED VIEW

Name

CREATE MATERIALIZED VIEW — define a new materialized view

Synopsis

CREATE MATERIALIZED VIEW [ IF NOT EXISTS ] table_name
[ (column_name [, ...] ) ]
[ WITH ( storage_parameter [= value] [, ... ] ) ]
[ TABLESPACE tablespace_name ]
AS query
[ WITH [ NO ] DATA ]

Description

CREATE MATERIALIZED VIEW defines a materialized view of a query. The query is executed and used to populate the view at the time the command is issued (unless WITH NO DATA is used) and may be refreshed later using REFRESH MATERIALIZED VIEW. CREATE MATERIALIZED VIEW is similar to CREATE TABLE AS, except that it also remembers the query used to initialize the view, so that it can be refreshed later upon demand. A materialized view has many of the same properties as a table, but there is no support for temporary materialized views or automatic generation of OIDs.

Parameters

IF NOT EXISTS

Do not throw an error if a materialized view with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing materialized view is anything like the one that would have been created.

table_name

The name (optionally schema-qualified) of the materialized view to be created.

column_name

The name of a column in the new materialized view. If column names are not provided, they are taken from the output column names of the query.

WITH ( storage_parameter [= value] [, ... ] )

This clause specifies optional storage parameters for the new materialized view; see Storage Parameters for more information. All parameters supported for CREATE TABLE are also supported for CREATE MATERIALIZED VIEW with the exception of OIDS. See CREATE TABLE for more information.

TABLESPACE tablespace_name

The tablespace_name is the name of the tablespace in which the new materialized view is to be created. If not specified, default_tablespace is consulted.
CREATE MATERIALIZED VIEW

query

A SELECT, TABLE, or VALUES command. This query will run within a security-restricted operation; in particular, calls to functions that themselves create temporary tables will fail.

WITH [ NO ] DATA

This clause specifies whether or not the materialized view should be populated at creation time. If not, the materialized view will be flagged as unscannable and cannot be queried until REFRESH MATERIALIZED VIEW is used.

Compatibility

CREATE MATERIALIZED VIEW is a PostgreSQL extension.

See Also

ALTER MATERIALIZED VIEW, CREATE TABLE AS, CREATE VIEW, DROP MATERIALIZED VIEW, REFRESH MATERIALIZED VIEW
CREATE OPERATOR

Name
CREATE OPERATOR — define a new operator

Synopsis
CREATE OPERATOR name {
    PROCEDURE = function_name
    [, LEFTARG = left_type ] [, RIGHTARG = right_type ]
    [, COMMUTATOR = com_op ] [, NEGATOR = neg_op ]
    [, RESTRICT = res_proc ] [, JOIN = join_proc ]
    [, HASHES ] [, MERGES ]
}

Description
CREATE OPERATOR defines a new operator, name. The user who defines an operator becomes its owner. If a schema name is given then the operator is created in the specified schema. Otherwise it is created in the current schema.

The operator name is a sequence of up to NAMEDATALEN-1 (63 by default) characters from the following list:

+ * / < > = ~ ! @ # % ^ & | ' ?

There are a few restrictions on your choice of name:

• -- and /* cannot appear anywhere in an operator name, since they will be taken as the start of a comment.
• A multicharacter operator name cannot end in + or -, unless the name also contains at least one of these characters:
  ~ ! @ # % ^ & | ' ?

For example, @- is an allowed operator name, but *- is not. This restriction allows PostgreSQL to parse SQL-compliant commands without requiring spaces between tokens.
• The use of -> as an operator name is deprecated. It may be disallowed altogether in a future release.

The operator != is mapped to <> on input, so these two names are always equivalent.

At least one of LEFTARG and RIGHTARG must be defined. For binary operators, both must be defined. For right unary operators, only LEFTARG should be defined, while for left unary operators only RIGHTARG should be defined.

The function_name procedure must have been previously defined using CREATE FUNCTION and must be defined to accept the correct number of arguments (either one or two) of the indicated types.

The other clauses specify optional operator optimization clauses. Their meaning is detailed in Section 36.13.
CREATE OPERATOR

To be able to create an operator, you must have USAGE privilege on the argument types and the return type, as well as EXECUTE privilege on the underlying function. If a commutator or negator operator is specified, you must own these operators.

Parameters

name
The name of the operator to be defined. See above for allowable characters. The name can be schema-qualified, for example CREATE OPERATOR myschema.+ (...). If not, then the operator is created in the current schema. Two operators in the same schema can have the same name if they operate on different data types. This is called overloading.

function_name
The function used to implement this operator.

left_type
The data type of the operator’s left operand, if any. This option would be omitted for a left-unary operator.

right_type
The data type of the operator’s right operand, if any. This option would be omitted for a right-unary operator.

com_op
The commutator of this operator.

neg_op
The negator of this operator.

res_proc
The restriction selectivity estimator function for this operator.

join_proc
The join selectivity estimator function for this operator.

HASES
Indicates this operator can support a hash join.

MERGES
Indicates this operator can support a merge join.

To give a schema-qualified operator name in com_op or the other optional arguments, use the OPERATOR() syntax, for example:

COMMUTATOR = OPERATOR(myschema.===) ,

Notes
Refer to Section 36.12 for further information.
CREATE OPERATOR

It is not possible to specify an operator’s lexical precedence in CREATE OPERATOR, because the parser’s precedence behavior is hard-wired. See Section 4.1.6 for precedence details.

The obsolete options SORT1, SORT2, LTCMP, and GTCMP were formerly used to specify the names of sort operators associated with a merge-joinable operator. This is no longer necessary, since information about associated operators is found by looking at B-tree operator families instead. If one of these options is given, it is ignored except for implicitly setting MERGES true.

Use DROP OPERATOR to delete user-defined operators from a database. Use ALTER OPERATOR to modify operators in a database.

Examples

The following command defines a new operator, area-equality, for the data type box:

```sql
CREATE OPERATOR === {
    LEFTARG = box,
    RIGHTARG = box,
    PROCEDURE = area_equal_procedure,
    COMMUTATOR = ===,
    NEGATOR = !==,
    RESTRICT = area_restriction_procedure,
    JOIN = area_join_procedure,
    HASHES, MERGES
};
```

Compatibility

CREATE OPERATOR is a PostgreSQL extension. There are no provisions for user-defined operators in the SQL standard.

See Also

ALTER OPERATOR, CREATE OPERATOR CLASS, DROP OPERATOR
CREATE OPERATOR CLASS

Name
CREATE OPERATOR CLASS — define a new operator class

Synopsis
CREATE OPERATOR CLASS name [ DEFAULT ] FOR TYPE data_type USING index_method [ FAMILY family_name ] AS
{ OPERATOR strategy_number operator_name [ ( op_type, op_type ) ] [ FOR SEARCH | FOR ORDER ]
 | FUNCTION support_number [ ( op_type [ , op_type ] ) ] function_name ( argument_type [, . . . ]
 | STORAGE storage_type
} [, . . . ]

Description
CREATE OPERATOR CLASS creates a new operator class. An operator class defines how a particular data type can be used with an index. The operator class specifies that certain operators will fill particular roles or “strategies” for this data type and this index method. The operator class also specifies the support procedures to be used by the index method when the operator class is selected for an index column. All the operators and functions used by an operator class must be defined before the operator class can be created.

If a schema name is given then the operator class is created in the specified schema. Otherwise it is created in the current schema. Two operator classes in the same schema can have the same name only if they are for different index methods.

The user who defines an operator class becomes its owner. Presently, the creating user must be a superuser. (This restriction is made because an erroneous operator class definition could confuse or even crash the server.)

CREATE OPERATOR CLASS does not presently check whether the operator class definition includes all the operators and functions required by the index method, nor whether the operators and functions form a self-consistent set. It is the user’s responsibility to define a valid operator class.

Related operator classes can be grouped into operator families. To add a new operator class to an existing family, specify the FAMILY option in CREATE OPERATOR CLASS. Without this option, the new class is placed into a family named the same as the new class (creating that family if it doesn’t already exist).

Refer to Section 36.14 for further information.

Parameters

name

The name of the operator class to be created. The name can be schema-qualified.

DEFAULT

If present, the operator class will become the default operator class for its data type. At most one operator class can be the default for a specific data type and index method.
CREATE OPERATOR CLASS

data_type
   The column data type that this operator class is for.

index_method
   The name of the index method this operator class is for.

family_name
   The name of the existing operator family to add this operator class to. If not specified, a family
   named the same as the operator class is used (creating it, if it doesn’t already exist).

strategy_number
   The index method’s strategy number for an operator associated with the operator class.

operator_name
   The name (optionally schema-qualified) of an operator associated with the operator class.

op_type
   In an OPERATOR clause, the operand data type(s) of the operator, or NONE to signify a left-unary
   or right-unary operator. The operand data types can be omitted in the normal case where they are
   the same as the operator class’s data type.

   In a FUNCTION clause, the operand data type(s) the function is intended to support, if different
   from the input data type(s) of the function (for B-tree comparison functions and hash functions)
   or the class’s data type (for B-tree sort support functions and all functions in GiST, SP-GiST, GIN
   and BRIN operator classes). These defaults are correct, and so op_type need not be specified in
   FUNCTION clauses, except for the case of a B-tree sort support function that is meant to support
   cross-data-type comparisons.

sort_family_name
   The name (optionally schema-qualified) of an existing btree operator family that describes the
   sort ordering associated with an ordering operator.

   If neither FOR SEARCH nor FOR ORDER BY is specified, FOR SEARCH is the default.

support_number
   The index method’s support procedure number for a function associated with the operator class.

function_name
   The name (optionally schema-qualified) of a function that is an index method support procedure
   for the operator class.

argument_type
   The parameter data type(s) of the function.

storage_type
   The data type actually stored in the index. Normally this is the same as the column data type,
   but some index methods (currently GiST, GIN and BRIN) allow it to be different. The STORAGE
   clause must be omitted unless the index method allows a different type to be used.

The OPERATOR, FUNCTION, and STORAGE clauses can appear in any order.
Notes

Because the index machinery does not check access permissions on functions before using them, including a function or operator in an operator class is tantamount to granting public execute permission on it. This is usually not an issue for the sorts of functions that are useful in an operator class.

The operators should not be defined by SQL functions. A SQL function is likely to be inlined into the calling query, which will prevent the optimizer from recognizing that the query matches an index.

Before PostgreSQL 8.4, the OPERATOR clause could include a RECHECK option. This is no longer supported because whether an index operator is “lossy” is now determined on-the-fly at run time. This allows efficient handling of cases where an operator might or might not be lossy.

Examples

The following example command defines a GiST index operator class for the data type _int4 (array of int4). See the intarray module for the complete example.

```
CREATE OPERATOR CLASS gist__int_ops
DEFAULT FOR TYPE _int4 USING gist AS
  OPERATOR 3 &&,
  OPERATOR 6 = (anyarray, anyarray),
  OPERATOR 7 >=,
  OPERATOR 8 <,
  OPERATOR 20 @@ (_int4, query_int),
  FUNCTION 1 g_int_consistent (internal, _int4, smallint, oid, internal),
  FUNCTION 2 g_int_union (internal, internal),
  FUNCTION 3 g_int_compress (internal),
  FUNCTION 4 g_int_decompress (internal),
  FUNCTION 5 g_int_penalty (internal, internal, internal),
  FUNCTION 6 g_int_picksplit (internal, internal),
  FUNCTION 7 g_int_same (_int4, _int4, internal);
```

Compatibility

CREATE OPERATOR CLASS is a PostgreSQL extension. There is no CREATE OPERATOR CLASS statement in the SQL standard.

See Also

ALTER OPERATOR CLASS, DROP OPERATOR CLASS, CREATE OPERATOR FAMILY, ALTER OPERATOR FAMILY
CREATE OPERATOR FAMILY

Name

CREATE OPERATOR FAMILY — define a new operator family

Synopsis

CREATE OPERATOR FAMILY name USING index_method

Description

CREATE OPERATOR FAMILY creates a new operator family. An operator family defines a collection of related operator classes, and perhaps some additional operators and support functions that are compatible with these operator classes but not essential for the functioning of any individual index. (Operators and functions that are essential to indexes should be grouped within the relevant operator class, rather than being "loose" in the operator family. Typically, single-data-type operators are bound to operator classes, while cross-data-type operators can be loose in an operator family containing operator classes for both data types.)

The new operator family is initially empty. It should be populated by issuing subsequent CREATE OPERATOR CLASS commands to add contained operator classes, and optionally ALTER OPERATOR FAMILY commands to add “loose” operators and their corresponding support functions.

If a schema name is given then the operator family is created in the specified schema. Otherwise it is created in the current schema. Two operator families in the same schema can have the same name only if they are for different index methods.

The user who defines an operator family becomes its owner. Presently, the creating user must be a superuser. (This restriction is made because an erroneous operator family definition could confuse or even crash the server.)

Refer to Section 36.14 for further information.

Parameters

name

The name of the operator family to be created. The name can be schema-qualified.

index_method

The name of the index method this operator family is for.

Compatibility

CREATE OPERATOR FAMILY is a PostgreSQL extension. There is no CREATE OPERATOR FAMILY statement in the SQL standard.
See Also

ALTER OPERATOR FAMILY, DROP OPERATOR FAMILY, CREATE OPERATOR CLASS, ALTER OPERATOR CLASS, DROP OPERATOR CLASS
CREATE POLICY

**Name**

CREATE POLICY — define a new row level security policy for a table

**Synopsis**

CREATE POLICY name ON table_name
[ FOR { ALL | SELECT | INSERT | UPDATE | DELETE } ]
[ TO { role_name | PUBLIC | CURRENT_USER | SESSION_USER } [, ...] ]
[ USING ( using_expression ) ]
[ WITH CHECK ( check_expression ) ]

**Description**

The CREATE POLICY command defines a new row-level security policy for a table. Note that row-level security must be enabled on the table (using ALTER TABLE ... ENABLE ROW LEVEL SECURITY) in order for created policies to be applied.

A policy grants the permission to select, insert, update, or delete rows that match the relevant policy expression. Existing table rows are checked against the expression specified in USING, while new rows that would be created via INSERT or UPDATE are checked against the expression specified in WITH CHECK. When a USING expression returns true for a given row then that row is visible to the user, while if false or null is returned then the row is not visible. When a WITH CHECK expression returns true for a row then that row is inserted or updated, while if false or null is returned then an error occurs.

For INSERT and UPDATE statements, WITH CHECK expressions are enforced after BEFORE triggers are fired, and before any actual data modifications are made. Thus a BEFORE ROW trigger may modify the data to be inserted, affecting the result of the security policy check. WITH CHECK expressions are enforced before any other constraints.

Policy names are per-table. Therefore, one policy name can be used for many different tables and have a definition for each table which is appropriate to that table.

Policies can be applied for specific commands or for specific roles. The default for newly created policies is that they apply for all commands and roles, unless otherwise specified. Multiple policies may apply to a single command; see below for more details. Table 1 summarizes how the different types of policy apply to specific commands.

For policies that can have both USING and WITH CHECK expressions (ALL and UPDATE), if no WITH CHECK expression is defined, then the USING expression will be used both to determine which rows are visible (normal USING case) and which new rows will be allowed to be added (WITH CHECK case).

If row-level security is enabled for a table, but no applicable policies exist, a “default deny” policy is assumed, so that no rows will be visible or updatable.
CREATE POLICY

Parameters

name
The name of the policy to be created. This must be distinct from the name of any other policy for the table.

table_name
The name (optionally schema-qualified) of the table the policy applies to.

command
The command to which the policy applies. Valid options are ALL, SELECT, INSERT, UPDATE, and DELETE. ALL is the default. See below for specifics regarding how these are applied.

role_name
The role(s) to which the policy is to be applied. The default is PUBLIC, which will apply the policy to all roles.

using_expression
Any SQL conditional expression (returning boolean). The conditional expression cannot contain any aggregate or window functions. This expression will be added to queries that refer to the table if row level security is enabled. Rows for which the expression returns true will be visible. Any rows for which the expression returns false or null will not be visible to the user (in a SELECT), and will not be available for modification (in an UPDATE or DELETE). Such rows are silently suppressed; no error is reported.

check_expression
Any SQL conditional expression (returning boolean). The conditional expression cannot contain any aggregate or window functions. This expression will be used in INSERT and UPDATE queries against the table if row level security is enabled. Only rows for which the expression evaluates to true will be allowed. An error will be thrown if the expression evaluates to false or null for any of the records inserted or any of the records that result from the update. Note that the check_expression is evaluated against the proposed new contents of the row, not the original contents.

Per-Command Policies

ALL
Using ALL for a policy means that it will apply to all commands, regardless of the type of command. If an ALL policy exists and more specific policies exist, then both the ALL policy and the more specific policy (or policies) will be applied. Additionally, ALL policies will be applied to both the selection side of a query and the modification side, using the USING expression for both cases if only a USING expression has been defined.

As an example, if an UPDATE is issued, then the ALL policy will be applicable both to what the UPDATE will be able to select as rows to be updated (applying the USING expression), and to the resulting updated rows, to check if they are permitted to be added to the table (applying the WITH CHECK expression, if defined, and the USING expression otherwise). If an INSERT or UPDATE command attempts to add rows to the table that do not pass the ALL policy’s WITH CHECK expression, the entire command will be aborted.
CREATE POLICY

SELECT

Using SELECT for a policy means that it will apply to SELECT queries and whenever SELECT permissions are required on the relation the policy is defined for. The result is that only those records from the relation that pass the SELECT policy will be returned during a SELECT query, and that queries that require SELECT permissions, such as UPDATE, will also only see those records that are allowed by the SELECT policy. A SELECT policy cannot have a WITH CHECK expression, as it only applies in cases where records are being retrieved from the relation.

INSERT

Using INSERT for a policy means that it will apply to INSERT commands. Rows being inserted that do not pass this policy will result in a policy violation error, and the entire INSERT command will be aborted. An INSERT policy cannot have a USING expression, as it only applies in cases where records are being added to the relation.

Note that INSERT with ON CONFLICT DO UPDATE checks INSERT policies’ WITH CHECK expressions only for rows appended to the relation by the INSERT path.

UPDATE

Using UPDATE for a policy means that it will apply to UPDATE, SELECT FOR UPDATE and SELECT FOR SHARE commands, as well as auxiliary ON CONFLICT DO UPDATE clauses of INSERT commands. Since UPDATE involves pulling an existing record and replacing it with a new modified record, UPDATE policies accept both a USING expression and a WITH CHECK expression. The USING expression determines which records the UPDATE command will see to operate against, while the WITH CHECK expression defines which modified rows are allowed to be stored back into the relation.

Any rows whose updated values do not pass the WITH CHECK expression will cause an error, and the entire command will be aborted. If only a USING clause is specified, then that clause will be used for both USING and WITH CHECK cases.

Typically an UPDATE command also needs to read data from columns in the relation being updated (e.g., in a WHERE clause or a RETURNING clause, or in an expression on the right hand side of the SET clause). In this case, SELECT rights are also required on the relation being updated, and the appropriate SELECT or ALL policies will be applied in addition to the UPDATE policies. Thus the user must have access to the row(s) being updated through a SELECT or ALL policy in addition to being granted permission to update the row(s) via an UPDATE or ALL policy.

When an INSERT command has an auxiliary ON CONFLICT DO UPDATE clause, if the UPDATE path is taken, the row to be updated is first checked against the USING expressions of any UPDATE policies, and then the new updated row is checked against the WITH CHECK expressions. Note, however, that unlike a standalone UPDATE command, if the existing row does not pass the USING expressions, an error will be thrown (the UPDATE path will never be silently avoided).

DELETE

Using DELETE for a policy means that it will apply to DELETE commands. Only rows that pass this policy will be seen by a DELETE command. There can be rows that are visible through a SELECT that are not available for deletion, if they do not pass the USING expression for the DELETE policy.

In most cases a DELETE command also needs to read data from columns in the relation that it is deleting from (e.g., in a WHERE clause or a RETURNING clause). In this case, SELECT rights are also required on the relation, and the appropriate SELECT or ALL policies will be applied in addition to the DELETE policies. Thus the user must have access to the row(s) being deleted through a SELECT or ALL policy in addition to being granted permission to delete the row(s) via a DELETE or ALL policy.
CREATE POLICY

A `DELETE` policy cannot have a `WITH CHECK` expression, as it only applies in cases where records are being deleted from the relation, so that there is no new row to check.

Table 1. Policies Applied by Command Type

<table>
<thead>
<tr>
<th>Command</th>
<th>SELECT/ALL policy</th>
<th>INSERT/ALL policy</th>
<th>UPDATE/ALL policy</th>
<th>DELETE/ALL policy</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>USING expression</td>
<td>WITH CHECK expression</td>
<td>USING expression</td>
<td>WITH CHECK expression</td>
</tr>
<tr>
<td>SELECT</td>
<td>Existing row</td>
<td>—</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>SELECT FOR UPDATE/SHARE</td>
<td>Existing row</td>
<td>—</td>
<td>Existing row</td>
<td>—</td>
</tr>
<tr>
<td>INSERT</td>
<td>—</td>
<td>New row</td>
<td>—</td>
<td>—</td>
</tr>
<tr>
<td>INSERT ... RETURNING</td>
<td>New row</td>
<td>—</td>
<td>New row</td>
<td>—</td>
</tr>
<tr>
<td>UPDATE</td>
<td>Existing &amp; new rows</td>
<td>—</td>
<td>Existing row</td>
<td>New row</td>
</tr>
<tr>
<td>DELETE</td>
<td>Existing row</td>
<td>—</td>
<td>—</td>
<td>Existing row</td>
</tr>
<tr>
<td>ON CONFLICT DO UPDATE</td>
<td>Existing &amp; new rows</td>
<td>—</td>
<td>Existing row</td>
<td>New row</td>
</tr>
</tbody>
</table>

Notes:

a. If read access is required to the existing or new row (for example, a `WHERE` or `RETURNING` clause that refers to columns from the relation).

Application of Multiple Policies

When multiple policies of different command types apply to the same command (for example, `SELECT` and `UPDATE` policies applied to an `UPDATE` command), then the user must have both types of permissions (for example, permission to select rows from the relation as well as permission to update them). Thus the expressions for one type of policy are combined with the expressions for the other type of policy using the `AND` operator.

When multiple policies of the same command type apply to the same command, then at least one of the policies must grant access to the relation. Thus the expressions from all the policies of that type are combined using the `OR` operator. If there are no applicable policies, then access is denied.

Note that, for the purposes of combining multiple policies, `ALL` policies are treated as having the same type as whichever other type of policy is being applied.

For example, in an `UPDATE` command requiring both `SELECT` and `UPDATE` permissions, if there are multiple applicable policies of each type, they will be combined as follows:

```
(expression from SELECT/ALL policy 1
 OR
 expression from SELECT/ALL policy 2
 OR
 ...)
 AND
```
CREATE POLICY

(expression from UPDATE/ALL policy 1
OR
expression from UPDATE/ALL policy 2
OR
...
)

Notes

You must be the owner of a table to create or change policies for it.

While policies will be applied for explicit queries against tables in the database, they are not applied when the system is performing internal referential integrity checks or validating constraints. This means there are indirect ways to determine that a given value exists. An example of this is attempting to insert a duplicate value into a column that is a primary key or has a unique constraint. If the insert fails then the user can infer that the value already exists. (This example assumes that the user is permitted by policy to insert records which they are not allowed to see.) Another example is where a user is allowed to insert into a table which references another, otherwise hidden table. Existence can be determined by the user inserting values into the referencing table, where success would indicate that the value exists in the referenced table. These issues can be addressed by carefully crafting policies to prevent users from being able to insert, delete, or update records at all which might possibly indicate a value they are not otherwise able to see, or by using generated values (e.g., surrogate keys) instead of keys with external meanings.

Generally, the system will enforce filter conditions imposed using security policies prior to qualifications that appear in user queries, in order to prevent inadvertent exposure of the protected data to user-defined functions which might not be trustworthy. However, functions and operators marked by the system (or the system administrator) as LEAKPROOF may be evaluated before policy expressions, as they are assumed to be trustworthy.

Since policy expressions are added to the user’s query directly, they will be run with the rights of the user running the overall query. Therefore, users who are using a given policy must be able to access any tables or functions referenced in the expression or they will simply receive a permission denied error when attempting to query the table that has row-level security enabled. This does not change how views work, however. As with normal queries and views, permission checks and policies for the tables which are referenced by a view will use the view owner’s rights and any policies which apply to the view owner.

Additional discussion and practical examples can be found in Section 5.7.

Compatibility

CREATE POLICY is a PostgreSQL extension.

See Also

ALTER POLICY, DROP POLICY, ALTER TABLE
CREATE ROLE

Name

CREATE ROLE — define a new database role

Synopsis

CREATE ROLE name [ [ WITH ] option [ ... ] ]

where option can be:

- SUPERUSER | NOSUPERUSER
- CREATEDB | NOCREATEDB
- CREATEROLE | NOCREATEROLE
- INHERIT | NOINHERIT
- LOGIN | NOLOGIN
- REPLICATION | NOREPLICATION
- BYPASSRLS | NOBYPASSRLS
- CONNECTION LIMIT connlimit
- [ ENCRYPTED | UNENCRYPTED ] PASSWORD ’password’
- VALID UNTIL ’timestamp’
- IN ROLE role_name [, ...]
- IN GROUP role_name [, ...]
- ROLE role_name [, ...]
- ADMIN role_name [, ...]
- USER role_name [, ...]
- SYSID uid

Description

CREATE ROLE adds a new role to a PostgreSQL database cluster. A role is an entity that can own database objects and have database privileges; a role can be considered a “user”, a “group”, or both depending on how it is used. Refer to Chapter 21 and Chapter 20 for information about managing users and authentication. You must have CREATEROLE privilege or be a database superuser to use this command.

Note that roles are defined at the database cluster level, and so are valid in all databases in the cluster.

Parameters

name

The name of the new role.

SUPERUSER
NOSUPERUSER

These clauses determine whether the new role is a “superuser”, who can override all access restrictions within the database. Superuser status is dangerous and should be used only when really needed. You must yourself be a superuser to create a new superuser. If not specified, NOSUPERUSER is the default.
CREATE ROLE

CREATEDB
NOCREATEDDB

These clauses define a role’s ability to create databases. If `CREATEDB` is specified, the role being defined will be allowed to create new databases. Specifying `NOCREATEDB` will deny a role the ability to create databases. If not specified, `NOCREATEDB` is the default.

CREATEROLE
NOCREATEROLE

These clauses determine whether a role will be permitted to create new roles (that is, execute `CREATE ROLE`). A role with `CREATEROLE` privilege can also alter and drop other roles. If not specified, `NOCREATEROLE` is the default.

INHERIT
NOINHERIT

These clauses determine whether a role “inherits” the privileges of roles it is a member of. A role with the `INHERIT` attribute can automatically use whatever database privileges have been granted to all roles it is directly or indirectly a member of. Without `INHERIT`, membership in another role only grants the ability to `SET ROLE` to that other role; the privileges of the other role are only available after having done so. If not specified, `INHERIT` is the default.

LOGIN
NOLOGIN

These clauses determine whether a role is allowed to log in; that is, whether the role can be given as the initial session authorization name during client connection. A role having the `LOGIN` attribute can be thought of as a user. Roles without this attribute are useful for managing database privileges, but are not users in the usual sense of the word. If not specified, `NOLOGIN` is the default, except when `CREATE ROLE` is invoked through its alternative spelling `CREATE USER`.

REPLICATION
NOREPLICATION

These clauses determine whether a role is a replication role. A role must have this attribute (or be a superuser) in order to be able to connect to the server in replication mode (physical or logical replication) and in order to be able to create or drop replication slots. A role having the `REPLICATION` attribute is a very highly privileged role, and should only be used on roles actually used for replication. If not specified, `NOREPLICATION` is the default.

BYPASSRLS
NOBYPASSRLS

These clauses determine whether a role bypasses every row-level security (RLS) policy. `NOBYPASSRLS` is the default. Note that `pg_dump` will set `row_security` to `OFF` by default, to ensure all contents of a table are dumped out. If the user running `pg_dump` does not have appropriate permissions, an error will be returned. The superuser and owner of the table being dumped always bypass RLS.

CONNECTION LIMIT connlimit

If role can log in, this specifies how many concurrent connections the role can make. -1 (the default) means no limit. Note that only normal connections are counted towards this limit. Neither prepared transactions nor background worker connections are counted towards this limit.

PASSWORD password

Sets the role’s password. (A password is only of use for roles having the `LOGIN` attribute, but you can nonetheless define one for roles without it.) If you do not plan to use password authentication you can omit this option. If no password is specified, the password will be set to null and
password authentication will always fail for that user. A null password can optionally be written explicitly as `PASSWORD NULL`.

**ENCRYPTED**  
**UNENCYPTED**  
These key words control whether the password is stored encrypted in the system catalogs. (If neither is specified, the default behavior is determined by the configuration parameter `password_encryption`.) If the presented password string is already in MD5-encrypted format, then it is stored encrypted as-is, regardless of whether `ENCRYPTED` or `UNENCYPTED` is specified (since the system cannot decrypt the specified encrypted password string). This allows reloading of encrypted passwords during dump/restore.

**VALID UNTIL 'timestamp'**  
The `VALID UNTIL` clause sets a date and time after which the role’s password is no longer valid. If this clause is omitted the password will be valid for all time.

**IN ROLE role_name**  
The `IN ROLE` clause lists one or more existing roles to which the new role will be immediately added as a new member. (Note that there is no option to add the new role as an administrator; use a separate `GRANT` command to do that.)

**IN GROUP role_name**  
`IN GROUP` is an obsolete spelling of `IN ROLE`.

**ROLE role_name**  
The `ROLE` clause lists one or more existing roles which are automatically added as members of the new role. (This in effect makes the new role a “group”.)

**ADMIN role_name**  
The `ADMIN` clause is like `ROLE`, but the named roles are added to the new role with `ADMIN OPTION`, giving them the right to grant membership in this role to others.

**USER role_name**  
The `USER` clause is an obsolete spelling of the `ROLE` clause.

**SYSID uid**  
The `SYSID` clause is ignored, but is accepted for backwards compatibility.

**Notes**

Use `ALTER ROLE` to change the attributes of a role, and `DROP ROLE` to remove a role. All the attributes specified by `CREATE ROLE` can be modified by later `ALTER ROLE` commands.

The preferred way to add and remove members of roles that are being used as groups is to use `GRANT` and `REVOKE`.

The `VALID UNTIL` clause defines an expiration time for a password only, not for the role per se. In particular, the expiration time is not enforced when logging in using a non-password-based authentication method.

The `INHERIT` attribute governs inheritance of grantable privileges (that is, access privileges for database objects and role memberships). It does not apply to the special role attributes set by `CREATE ROLE` and `ALTER ROLE`. For example, being a member of a role with `CREATEDB` privilege does not
immediately grant the ability to create databases, even if INHERIT is set; it would be necessary to become that role via SET ROLE before creating a database.

The INHERIT attribute is the default for reasons of backwards compatibility: in prior releases of PostgreSQL, users always had access to all privileges of groups they were members of. However, NOINHERIT provides a closer match to the semantics specified in the SQL standard.

Be careful with the CREATEROLE privilege. There is no concept of inheritance for the privileges of a CREATEROLE-role. That means that even if a role does not have a certain privilege but is allowed to create other roles, it can easily create another role with different privileges than its own (except for creating roles with superuser privileges). For example, if the role “user” has the CREATEROLE privilege but not the CREATEDB privilege, nonetheless it can create a new role with the CREATEDB privilege. Therefore, regard roles that have the CREATEROLE privilege as almost-superuser-roles.

PostgreSQL includes a program createuser that has the same functionality as CREATE ROLE (in fact, it calls this command) but can be run from the command shell.

The CONNECTION LIMIT option is only enforced approximately; if two new sessions start at about the same time when just one connection “slot” remains for the role, it is possible that both will fail. Also, the limit is never enforced for superusers.

Caution must be exercised when specifying an unencrypted password with this command. The password will be transmitted to the server in cleartext, and it might also be logged in the client’s command history or the server log. The command createuser, however, transmits the password encrypted. Also, psql contains a command \password that can be used to safely change the password later.

Examples

Create a role that can log in, but don’t give it a password:

```
CREATE ROLE jonathan LOGIN;
```

Create a role with a password:

```
CREATE USER davide WITH PASSWORD 'jw8s0F4';
```

(create USER is the same as CREATE ROLE except that it implies LOGIN.)

Create a role with a password that is valid until the end of 2004. After one second has ticked in 2005, the password is no longer valid.

```
CREATE ROLE miriam WITH LOGIN PASSWORD 'jw8s0F4' VALID UNTIL '2005-01-01';
```

Create a role that can create databases and manage roles:

```
CREATE ROLE admin WITH CREATEDB CREATEROLE;
```
**Compatibility**

The `CREATE ROLE` statement is in the SQL standard, but the standard only requires the syntax

```
CREATE ROLE name [ WITH ADMIN role_name ]
```

Multiple initial administrators, and all the other options of `CREATE ROLE`, are PostgreSQL extensions. The SQL standard defines the concepts of users and roles, but it regards them as distinct concepts and leaves all commands defining users to be specified by each database implementation. In PostgreSQL we have chosen to unify users and roles into a single kind of entity. Roles therefore have many more optional attributes than they do in the standard.

The behavior specified by the SQL standard is most closely approximated by giving users the `NOINHERIT` attribute, while roles are given the `INHERIT` attribute.

**See Also**

SET ROLE, ALTER ROLE, DROP ROLE, GRANT, REVOKE, createuser
CREATE RULE

Name
CREATE RULE — define a new rewrite rule

Synopsis
CREATE [ OR REPLACE ] RULE name AS ON event
    TO table_name [ WHERE condition ]
    DO [ ALSO | INSTEAD ] { NOTHING | command | (command; command...) }

where event can be one of:

    SELECT | INSERT | UPDATE | DELETE

Description
CREATE RULE defines a new rule applying to a specified table or view. CREATE OR REPLACE RULE will either create a new rule, or replace an existing rule of the same name for the same table.

The PostgreSQL rule system allows one to define an alternative action to be performed on insertions, updates, or deletions in database tables. Roughly speaking, a rule causes additional commands to be executed when a given command on a given table is executed. Alternatively, an INSTEAD rule can replace a given command by another, or cause a command not to be executed at all. Rules are used to implement SQL views as well. It is important to realize that a rule is really a command transformation mechanism, or command macro. The transformation happens before the execution of the command starts. If you actually want an operation that fires independently for each physical row, you probably want to use a trigger, not a rule. More information about the rules system is in Chapter 39.

Presently, ON SELECT rules must be unconditional INSTEAD rules and must have actions that consist of a single SELECT command. Thus, an ON SELECT rule effectively turns the table into a view, whose visible contents are the rows returned by the rule’s SELECT command rather than whatever had been stored in the table (if anything). It is considered better style to write a CREATE VIEW command than to create a real table and define an ON SELECT rule for it.

You can create the illusion of an updatable view by defining ON INSERT, ON UPDATE, and ON DELETE rules (or any subset of those that’s sufficient for your purposes) to replace update actions on the view with appropriate updates on other tables. If you want to support INSERT RETURNING and so on, then be sure to put a suitable RETURNING clause into each of these rules.

There is a catch if you try to use conditional rules for complex view updates: there must be an unconditional INSTEAD rule for each action you wish to allow on the view. If the rule is conditional, or is not INSTEAD, then the system will still reject attempts to perform the update action, because it thinks it might end up trying to perform the action on the dummy table of the view in some cases. If you want to handle all the useful cases in conditional rules, add an unconditional DO INSTEAD NOTHING rule to ensure that the system understands it will never be called on to update the dummy table. Then make the conditional rules non-INSTEAD; in the cases where they are applied, they add to the default INSTEAD NOTHING action. (This method does not currently work to support RETURNING queries, however.)
Note: A view that is simple enough to be automatically updatable (see CREATE VIEW) does not require a user-created rule in order to be updatable. While you can create an explicit rule anyway, the automatic update transformation will generally outperform an explicit rule.

Another alternative worth considering is to use INSTEAD OF triggers (see CREATE TRIGGER) in place of rules.

Parameters

name

The name of a rule to create. This must be distinct from the name of any other rule for the same table. Multiple rules on the same table and same event type are applied in alphabetical name order.

event

The event is one of SELECT, INSERT, UPDATE, or DELETE. Note that an INSERT containing an ON CONFLICT clause cannot be used on tables that have either INSERT or UPDATE rules. Consider using an updatable view instead.

table_name

The name (optionally schema-qualified) of the table or view the rule applies to.

condition

Any SQL conditional expression (returning boolean). The condition expression cannot refer to any tables except NEW and OLD, and cannot contain aggregate functions.

INSTEAD

INSTEAD indicates that the commands should be executed instead of the original command.

ALSO

ALSO indicates that the commands should be executed in addition to the original command.

If neither ALSO nor INSTEAD is specified, ALSO is the default.

command

The command or commands that make up the rule action. Valid commands are SELECT, INSERT, UPDATE, DELETE, or NOTIFY.

Within condition and command, the special table names NEW and OLD can be used to refer to values in the referenced table. NEW is valid in ON INSERT and ON UPDATE rules to refer to the new row being inserted or updated. OLD is valid in ON UPDATE and ON DELETE rules to refer to the existing row being updated or deleted.

Notes

You must be the owner of a table to create or change rules for it.

In a rule for INSERT, UPDATE, or DELETE on a view, you can add a RETURNING clause that emits the view’s columns. This clause will be used to compute the outputs if the rule is triggered by an INSERT RETURNING, UPDATE RETURNING, or DELETE RETURNING command respectively. When the rule is triggered by a command without RETURNING, the rule’s RETURNING clause will be ignored. The
current implementation allows only unconditional \texttt{INSTEAD} rules to contain \texttt{RETURNING}; furthermore there can be at most one \texttt{RETURNING} clause among all the rules for the same event. (This ensures that there is only one candidate \texttt{RETURNING} clause to be used to compute the results.) \texttt{RETURNING} queries on the view will be rejected if there is no \texttt{RETURNING} clause in any available rule.

It is very important to take care to avoid circular rules. For example, though each of the following two rule definitions are accepted by PostgreSQL, the \texttt{SELECT} command would cause PostgreSQL to report an error because of recursive expansion of a rule:

\begin{verbatim}
CREATE RULE "_RETURN" AS
  ON SELECT TO t1
  DO INSTEAD
    SELECT * FROM t2;

CREATE RULE "_RETURN" AS
  ON SELECT TO t2
  DO INSTEAD
    SELECT * FROM t1;

SELECT * FROM t1;
\end{verbatim}

Presently, if a rule action contains a \texttt{NOTIFY} command, the \texttt{NOTIFY} command will be executed unconditionally, that is, the \texttt{NOTIFY} will be issued even if there are not any rows that the rule should apply to. For example, in:

\begin{verbatim}
CREATE RULE notify_me AS ON UPDATE TO mytable DO ALSO NOTIFY mytable;

UPDATE mytable SET name = 'foo' WHERE id = 42;
\end{verbatim}

one \texttt{NOTIFY} event will be sent during the \texttt{UPDATE}, whether or not there are any rows that match the condition \texttt{id = 42}. This is an implementation restriction that might be fixed in future releases.

\section*{Compatibility}

\texttt{CREATE RULE} is a PostgreSQL language extension, as is the entire query rewrite system.

\section*{See Also}

\texttt{ALTER RULE, DROP RULE}
CREATE SCHEMA

Name
CREATE SCHEMA — define a new schema

Synopsis
CREATE SCHEMA schema_name [ AUTHORIZATION role_specification ] [ schema_element [ ... ] ]
CREATE SCHEMA AUTHORIZATION role_specification [ schema_element [ ... ] ]
CREATE SCHEMA IF NOT EXISTS schema_name [ AUTHORIZATION role_specification ]
CREATE SCHEMA IF NOT EXISTS AUTHORIZATION role_specification

where role_specification can be:

user_name
| CURRENT_USER
| SESSION_USER

Description
CREATE SCHEMA enters a new schema into the current database. The schema name must be distinct from the name of any existing schema in the current database.

A schema is essentially a namespace: it contains named objects (tables, data types, functions, and operators) whose names can duplicate those of other objects existing in other schemas. Named objects are accessed either by “qualifying” their names with the schema name as a prefix, or by setting a search path that includes the desired schema(s). A CREATE command specifying an unqualified object name creates the object in the current schema (the one at the front of the search path, which can be determined with the function current_schema).

Optionally, CREATE SCHEMA can include subcommands to create objects within the new schema. The subcommands are treated essentially the same as separate commands issued after creating the schema, except that if the AUTHORIZATION clause is used, all the created objects will be owned by that user.

Parameters

schema_name

The name of a schema to be created. If this is omitted, the user_name is used as the schema name. The name cannot begin with pg_, as such names are reserved for system schemas.

user_name

The role name of the user who will own the new schema. If omitted, defaults to the user executing the command. To create a schema owned by another role, you must be a direct or indirect member of that role, or be a superuser.

schema_element

An SQL statement defining an object to be created within the schema. Currently, only CREATE TABLE, CREATE VIEW, CREATE INDEX, CREATE SEQUENCE, CREATE TRIGGER and GRANT
CREATE SCHEMA

are accepted as clauses within CREATE SCHEMA. Other kinds of objects may be created in separate commands after the schema is created.

IF NOT EXISTS

Do nothing (except issuing a notice) if a schema with the same name already exists. schema_element subcommands cannot be included when this option is used.

Notes

To create a schema, the invoking user must have the CREATE privilege for the current database. (Of course, superusers bypass this check.)

Examples

Create a schema:

CREATE SCHEMA myschema;

Create a schema for user joe; the schema will also be named joe:

CREATE SCHEMA AUTHORIZATION joe;

Create a schema named test that will be owned by user joe, unless there already is a schema named test. (It does not matter whether joe owns the pre-existing schema.)

CREATE SCHEMA IF NOT EXISTS test AUTHORIZATION joe;

Create a schema and create a table and view within it:

CREATE SCHEMA hollywood
  CREATE TABLE films (title text, release date, awards text[])
  CREATE VIEW winners AS
    SELECT title, release FROM films WHERE awards IS NOT NULL;

Notice that the individual subcommands do not end with semicolons.

The following is an equivalent way of accomplishing the same result:

CREATE SCHEMA hollywood;
CREATE TABLE hollywood.films (title text, release date, awards text[]);
CREATE VIEW hollywood.winners AS
  SELECT title, release FROM hollywood.films WHERE awards IS NOT NULL;


**Compatibility**

The SQL standard allows a `DEFAULT CHARACTER SET` clause in `CREATE SCHEMA`, as well as more subcommand types than are presently accepted by PostgreSQL.

The SQL standard specifies that the subcommands in `CREATE SCHEMA` can appear in any order. The present PostgreSQL implementation does not handle all cases of forward references in subcommands; it might sometimes be necessary to reorder the subcommands in order to avoid forward references.

According to the SQL standard, the owner of a schema always owns all objects within it. PostgreSQL allows schemas to contain objects owned by users other than the schema owner. This can happen only if the schema owner grants the `CREATE` privilege on their schema to someone else, or a superuser chooses to create objects in it.

The **IF NOT EXISTS** option is a PostgreSQL extension.

**See Also**

ALTER SCHEMA, DROP SCHEMA
CREATE SEQUENCE

**Name**

CREATE SEQUENCE — define a new sequence generator

**Synopsis**

CREATE [ TEMPORARY | TEMP | SEQUENCE [ IF NOT EXISTS ] name [ INCREMENT [ BY ] increment ]
 [ MINVALUE minvalue | NO MINVALUE ] [ MAXVALUE maxvalue | NO MAXVALUE ]
 [ START [ WITH ] start ] [ CACHE cache ] [ [ NO ] CYCLE ]
 [ OWNED BY { table_name.column_name | NONE } ]

**Description**

CREATE SEQUENCE creates a new sequence number generator. This involves creating and initializing a new special single-row table with the name *name*. The generator will be owned by the user issuing the command.

If a schema name is given then the sequence is created in the specified schema. Otherwise it is created in the current schema. Temporary sequences exist in a special schema, so a schema name cannot be given when creating a temporary sequence. The sequence name must be distinct from the name of any other sequence, table, index, view, or foreign table in the same schema.

After a sequence is created, you use the functions `nextval`, `currval`, and `setval` to operate on the sequence. These functions are documented in Section 9.16.

Although you cannot update a sequence directly, you can use a query like:

```
SELECT * FROM name;
```

to examine the parameters and current state of a sequence. In particular, the `last_value` field of the sequence shows the last value allocated by any session. (Of course, this value might be obsolete by the time it’s printed, if other sessions are actively doing `nextval` calls.)

**Parameters**

**TEMPORARY** or **TEMP**

If specified, the sequence object is created only for this session, and is automatically dropped on session exit. Existing permanent sequences with the same name are not visible (in this session) while the temporary sequence exists, unless they are referenced with schema-qualified names.

**IF NOT EXISTS**

Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing relation is anything like the sequence that would have been created - it might not even be a sequence.

**name**

The name (optionally schema-qualified) of the sequence to be created.
CREATE SEQUENCE

increment

The optional clause INCREMENT BY increment specifies which value is added to the current sequence value to create a new value. A positive value will make an ascending sequence, a negative one a descending sequence. The default value is 1.

minvalue

NO MINVALUE

The optional clause MINVALUE minvalue determines the minimum value a sequence can generate. If this clause is not supplied or NO MINVALUE is specified, then defaults will be used. The defaults are 1 and \(-2^{64} - 1\) for ascending and descending sequences, respectively.

maxvalue

NO MAXVALUE

The optional clause MAXVALUE maxvalue determines the maximum value for the sequence. If this clause is not supplied or NO MAXVALUE is specified, then default values will be used. The defaults are \(2^{63} - 1\) and \(-1\) for ascending and descending sequences, respectively.

start

The optional clause START WITH start allows the sequence to begin anywhere. The default starting value is minvalue for ascending sequences and maxvalue for descending ones.

cache

The optional clause CACHE cache specifies how many sequence numbers are to be preallocated and stored in memory for faster access. The minimum value is 1 (only one value can be generated at a time, i.e., no cache), and this is also the default.

CYCLE

NO CYCLE

The CYCLE option allows the sequence to wrap around when the maxvalue or minvalue has been reached by an ascending or descending sequence respectively. If the limit is reached, the next number generated will be the minvalue or maxvalue, respectively.

If NO CYCLE is specified, any calls to nextval after the sequence has reached its maximum value will return an error. If neither CYCLE or NO CYCLE are specified, NO CYCLE is the default.

OWNED BY table_name.column_name

OWNED BY NONE

The OWNED BY option causes the sequence to be associated with a specific table column, such that if that column (or its whole table) is dropped, the sequence will be automatically dropped as well. The specified table must have the same owner and be in the same schema as the sequence. OWNED BY NONE, the default, specifies that there is no such association.

Notes

Use DROP SEQUENCE to remove a sequence.

Sequences are based on bigint arithmetic, so the range cannot exceed the range of an eight-byte integer (-9223372036854775808 to 9223372036854775807).

Because nextval and setval calls are never rolled back, sequence objects cannot be used if “gapless” assignment of sequence numbers is needed. It is possible to build gapless assignment by using exclusive locking of a table containing a counter; but this solution is much more expensive than sequence objects, especially if many transactions need sequence numbers concurrently.
Unexpected results might be obtained if a cache setting greater than one is used for a sequence object that will be used concurrently by multiple sessions. Each session will allocate and cache successive sequence values during one access to the sequence object and increase the sequence object’s last_value accordingly. Then, the next cache-1 uses of nextval within that session simply return the preallocated values without touching the sequence object. So, any numbers allocated but not used within a session will be lost when that session ends, resulting in “holes” in the sequence.

Furthermore, although multiple sessions are guaranteed to allocate distinct sequence values, the values might be generated out of sequence when all the sessions are considered. For example, with a cache setting of 10, session A might reserve values 1..10 and return nextval=1, then session B might reserve values 11..20 and return nextval=11 before session A has generated nextval=2. Thus, with a cache setting of one it is safe to assume that nextval values are generated sequentially; with a cache setting greater than one you should only assume that the nextval values are all distinct, not that they are generated purely sequentially. Also, last_value will reflect the latest value reserved by any session, whether or not it has yet been returned by nextval.

Another consideration is that a setval executed on such a sequence will not be noticed by other sessions until they have used up any preallocated values they have cached.

**Examples**

Create an ascending sequence called serial, starting at 101:

```sql
CREATE SEQUENCE serial START 101;
```

Select the next number from this sequence:

```sql
SELECT nextval('serial');
```

```
nextval
-------
101
```

Select the next number from this sequence:

```sql
SELECT nextval('serial');
```

```
nextval
-------
102
```

Use this sequence in an INSERT command:

```sql
INSERT INTO distributors VALUES (nextval('serial'), 'nothing');
```

Update the sequence value after a COPY FROM:

```sql
BEGIN;
COPY distributors FROM 'input_file';
SELECT setval('serial', max(id)) FROM distributors;
```
CREATE SEQUENCE

END;

Compatibility

CREATE SEQUENCE conforms to the SQL standard, with the following exceptions:

- The standard’s AS data_type expression is not supported.
- Obtaining the next value is done using the nextval() function instead of the standard’s NEXT VALUE FOR expression.
- The OWNED BY clause is a PostgreSQL extension.

See Also

ALTER SEQUENCE, DROP SEQUENCE
CREATE SERVER

Name
CREATE SERVER — define a new foreign server

Synopsis
CREATE SERVER server_name [ TYPE 'server_type' ] [ VERSION 'server_version' ]
FOREIGN DATA WRAPPER fdw_name
[ OPTIONS ( option 'value' [, ... ] ) ]

Description
CREATE SERVER defines a new foreign server. The user who defines the server becomes its owner.
A foreign server typically encapsulates connection information that a foreign-data wrapper uses to
access an external data resource. Additional user-specific connection information may be specified by
means of user mappings.
The server name must be unique within the database.
Creating a server requires USAGE privilege on the foreign-data wrapper being used.

Parameters

server_name
    The name of the foreign server to be created.
server_type
    Optional server type, potentially useful to foreign-data wrappers.
server_version
    Optional server version, potentially useful to foreign-data wrappers.
fdw_name
    The name of the foreign-data wrapper that manages the server.
OPTIONS ( option 'value' [, ... ] )
    This clause specifies the options for the server. The options typically define the connection details
    of the server, but the actual names and values are dependent on the server’s foreign-data wrapper.

Notes
When using the dblink module, a foreign server’s name can be used as an argument of the
dblink_connect function to indicate the connection parameters. It is necessary to have the USAGE
privilege on the foreign server to be able to use it in this way.
**CREATE SERVER**

**Examples**

Create a server `myserver` that uses the foreign-data wrapper `postgres_fdw`:

```
CREATE SERVER myserver FOREIGN DATA WRAPPER postgres_fdw OPTIONS (host 'foo', dbname 'foodb', port '5432');
```

See `postgres_fdw` for more details.

**Compatibility**

`CREATE SERVER` conforms to ISO/IEC 9075-9 (SQL/MED).

**See Also**

`ALTER SERVER`, `DROP SERVER`, `CREATE FOREIGN DATA WRAPPER`, `CREATE FOREIGN TABLE`, `CREATE USER MAPPING`
CREATE TABLE

Name
CREATE TABLE — define a new table

Synopsis

CREATE [ [ GLOBAL | LOCAL ] { TEMPORARY | TEMP } | UNLOGGED ] TABLE [ IF NOT EXISTS ] table_name
{ column_name data_type [ COLLATE collation ] [ column_constraint [ ... ] ]
| table_constraint
| LIKE source_table [ like_option ... ] }
[, ... ]
} [ INHERITS ( parent_table [, ... ] ) ]
[ WITH { storage_parameter [ = value ] [, ... ] } | WITH OIDS | WITHOUT OIDS ]
[ ON COMMIT { PRESERVE ROWS | DELETE ROWS | DROP } ]
[ TABLESPACE tablespace_name ]

CREATE [ [ GLOBAL | LOCAL ] { TEMPORARY | TEMP } | UNLOGGED ] TABLE [ IF NOT EXISTS ] table_name
OF type_name [ (
{ column_name WITH OPTIONS [ column_constraint [ ... ] ]
| table_constraint }
[, ... ]
) ] [ WITH { storage_parameter [ = value ] [, ... ] } | WITH OIDS | WITHOUT OIDS ]
[ ON COMMIT { PRESERVE ROWS | DELETE ROWS | DROP } ]
[ TABLESPACE tablespace_name ]

where column_constraint is:

{ CONSTRAINT constraint_name ]
| NOT NULL |
| NULL |
| CHECK ( expression ) [ NO INHERIT ] |
| DEFAULT default_expr |
| UNIQUE index_parameters |
| PRIMARY KEY index_parameters |
| REFERENCES reftable [ ( refcolumn ) ] [ MATCH FULL | MATCH PARTIAL | MATCH SIMPLE ]
| [ ON DELETE action ] [ ON UPDATE action ] ]
| DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ]

and table_constraint is:

{ CONSTRAINT constraint_name ]
| CHECK ( expression ) [ NO INHERIT ] |
| UNIQUE ( column_name [, ... ] ) index_parameters |
| PRIMARY KEY ( column_name [, ... ] ) index_parameters |
| EXCLUDE [ USING index_method ] ( exclude_element WITH operator [, ... ] ) index_parameters |
| FOREIGN KEY ( column_name [, ... ] ) REFERENCES reftable [ ( refcolumn [, ... ] ) ]
| [ MATCH FULL | MATCH PARTIAL | MATCH SIMPLE ] [ ON DELETE action ] [ ON UPDATE action ]
| DEFERRABLE | NOT DEFERRABLE ] [ INITIALLY DEFERRED | INITIALLY IMMEDIATE ]

and like_option is:

1508
CREATE TABLE

{ INCLUDING | EXCLUDING } { DEFAULTS | CONSTRAINTS | INDEXES | STORAGE | COMMENTS | ALL }

index_parameters in UNIQUE, PRIMARY KEY, and EXCLUDE constraints are:

[ WITH ( storage_parameter [= value] [, ... ] ) ]
[ USING INDEX TABLESPACE tablespace_name ]

exclude_element in an EXCLUDE constraint is:

{ column_name | ( expression ) } [ opclass ] [ ASC | DESC ] [ NULLS { FIRST | LAST } ]

Description

CREATE TABLE will create a new, initially empty table in the current database. The table will be owned by the user issuing the command.

If a schema name is given (for example, CREATE TABLE myschema.mytable ...) then the table is created in the specified schema. Otherwise it is created in the current schema. Temporary tables exist in a special schema, so a schema name cannot be given when creating a temporary table. The name of the table must be distinct from the name of any other table, sequence, index, view, or foreign table in the same schema.

CREATE TABLE also automatically creates a data type that represents the composite type corresponding to one row of the table. Therefore, tables cannot have the same name as any existing data type in the same schema.

The optional constraint clauses specify constraints (tests) that new or updated rows must satisfy for an insert or update operation to succeed. A constraint is an SQL object that helps define the set of valid values in the table in various ways.

There are two ways to define constraints: table constraints and column constraints. A column constraint is defined as part of a column definition. A table constraint definition is not tied to a particular column, and it can encompass more than one column. Every column constraint can also be written as a table constraint; a column constraint is only a notational convenience for use when the constraint only affects one column.

To be able to create a table, you must have USAGE privilege on all column types or the type in the OF clause, respectively.

Parameters

TEMPORARY or TEMP

If specified, the table is created as a temporary table. Temporary tables are automatically dropped at the end of a session, or optionally at the end of the current transaction (see ON COMMIT below). Existing permanent tables with the same name are not visible to the current session while the temporary table exists, unless they are referenced with schema-qualified names. Any indexes created on a temporary table are automatically temporary as well.

The autovacuum daemon cannot access and therefore cannot vacuum or analyze temporary tables. For this reason, appropriate vacuum and analyze operations should be performed via session SQL commands. For example, if a temporary table is going to be used in complex queries, it is wise to run ANALYZE on the temporary table after it is populated.
CREATE TABLE

Optionally, \texttt{GLOBAL} or \texttt{LOCAL} can be written before \texttt{TEMPORARY} or \texttt{TEMP}. This presently makes no difference in PostgreSQL and is deprecated; see \textit{Compatibility}.

\textbf{UNLOGGED}

If specified, the table is created as an unlogged table. Data written to unlogged tables is not written to the write-ahead log (see Chapter 30), which makes them considerably faster than ordinary tables. However, they are not crash-safe: an unlogged table is automatically truncated after a crash or unclean shutdown. The contents of an unlogged table are also not replicated to standby servers. Any indexes created on an unlogged table are automatically unlogged as well.

\textbf{IF NOT EXISTS}

Do not throw an error if a relation with the same name already exists. A notice is issued in this case. Note that there is no guarantee that the existing relation is anything like the one that would have been created.

\texttt{table\_name}

The name (optionally schema-qualified) of the table to be created.

\textbf{OF type\_name}

Creates a \textit{typed table}, which takes its structure from the specified composite type (name optionally schema-qualified). A typed table is tied to its type; for example the table will be dropped if the type is dropped (with \texttt{DROP TYPE \ldots CASCADE}).

When a typed table is created, then the data types of the columns are determined by the underlying composite type and are not specified by the \texttt{CREATE TABLE} command. But the \texttt{CREATE TABLE} command can add defaults and constraints to the table and can specify storage parameters.

\texttt{column\_name}

The name of a column to be created in the new table.

\texttt{data\_type}

The data type of the column. This can include array specifiers. For more information on the data types supported by PostgreSQL, refer to Chapter 8.

\textbf{COLLATE collation}

The \texttt{COLLATE} clause assigns a collation to the column (which must be of a collatable data type). If not specified, the column data type’s default collation is used.

\textbf{INHERITS ( parent\_table [, \ldots ] )}

The optional \texttt{INHERITS} clause specifies a list of tables from which the new table automatically inherits all columns. Parent tables can be plain tables or foreign tables.

Use of \texttt{INHERITS} creates a persistent relationship between the new child table and its parent table(s). Schema modifications to the parent(s) normally propagate to children as well, and by default the data of the child table is included in scans of the parent(s).

If the same column name exists in more than one parent table, an error is reported unless the data types of the columns match in each of the parent tables. If there is no conflict, then the duplicate columns are merged to form a single column in the new table. If the column name list of the new table contains a column name that is also inherited, the data type must likewise match the inherited column(s), and the column definitions are merged into one. If the new table explicitly specifies a default value for the column, this default overrides any defaults from inherited declarations of the column. Otherwise, any parents that specify default values for the column must all specify the same default, or an error will be reported.
CREATE TABLE

CHECK constraints are merged in essentially the same way as columns: if multiple parent tables and/or the new table definition contain identically-named CHECK constraints, these constraints must all have the same check expression, or an error will be reported. Constraints having the same name and expression will be merged into one copy. A constraint marked NO INHERIT in a parent will not be considered. Notice that an unnamed CHECK constraint in the new table will never be merged, since a unique name will always be chosen for it.

Column STORAGE settings are also copied from parent tables.

LIKE source_table [ like_option ... ]

The LIKE clause specifies a table from which the new table automatically copies all column names, their data types, and their not-null constraints.

Unlike INHERITS, the new table and original table are completely decoupled after creation is complete. Changes to the original table will not be applied to the new table, and it is not possible to include data of the new table in scans of the original table.

Default expressions for the copied column definitions will be copied only if INCLUDING DEFAULTS is specified. The default behavior is to exclude default expressions, resulting in the copied columns in the new table having null defaults. Note that copying defaults that call database-modification functions, such as nextval, may create a functional linkage between the original and new tables.

Not-null constraints are always copied to the new table. CHECK constraints will be copied only if INCLUDING CONSTRAINTS is specified. No distinction is made between column constraints and table constraints.

Indexes, PRIMARY KEY, UNIQUE, and EXCLUDE constraints on the original table will be created on the new table only if INCLUDING INDEXES is specified. Names for the new indexes and constraints are chosen according to the default rules, regardless of how the originals were named. (This behavior avoids possible duplicate-name failures for the new indexes.)

STORAGE settings for the copied column definitions will be copied only if INCLUDING STORAGE is specified. The default behavior is to exclude STORAGE settings, resulting in the copied columns in the new table having type-specific default settings. For more on STORAGE settings, see Section 65.2.

Comments for the copied columns, constraints, and indexes will be copied only if INCLUDING COMMENTS is specified. The default behavior is to exclude comments, resulting in the copied columns and constraints in the new table having no comments.

INCLUDING ALL is an abbreviated form of INCLUDING DEFAULTS INCLUDING CONSTRAINTS INCLUDING INDEXES INCLUDING STORAGE INCLUDING COMMENTS.

Note that unlike INHERITS, columns and constraints copied by LIKE are not merged with similarly named columns and constraints. If the same name is specified explicitly or in another LIKE clause, an error is signaled.

The LIKE clause can also be used to copy column definitions from views, foreign tables, or composite types. Inapplicable options (e.g., INCLUDING INDEXES from a view) are ignored.

CONSTRAINT constraint_name

An optional name for a column or table constraint. If the constraint is violated, the constraint name is present in error messages, so constraint names like col must be positive can be used to communicate helpful constraint information to client applications. (Double-quotes are needed to specify constraint names that contain spaces.) If a constraint name is not specified, the system generates a name.
NOT NULL

The column is not allowed to contain null values.

NULL

The column is allowed to contain null values. This is the default.

This clause is only provided for compatibility with non-standard SQL databases. Its use is dis-couraged in new applications.

CHECK ( expression ) [ NO INHERIT ]

The CHECK clause specifies an expression producing a Boolean result which new or updated rows must satisfy for an insert or update operation to succeed. Expressions evaluating to TRUE or UNKNOWN succeed. Should any row of an insert or update operation produce a FALSE result, an error exception is raised and the insert or update does not alter the database. A check constraint specified as a column constraint should reference that column’s value only, while an expression appearing in a table constraint can reference multiple columns.

Currently, CHECK expressions cannot contain subqueries nor refer to variables other than columns of the current row (see Section 5.3.1). The system column tableoid may be referenced, but not any other system column.

A constraint marked with NO INHERIT will not propagate to child tables.

When a table has multiple CHECK constraints, they will be tested for each row in alphabetical order by name, after checking NOT NULL constraints. (PostgreSQL versions before 9.5 did not honor any particular firing order for CHECK constraints.)

DEFAULT default_expr

The DEFAULT clause assigns a default data value for the column whose column definition it appears within. The value is any variable-free expression (subqueries and cross-references to other columns in the current table are not allowed). The data type of the default expression must match the data type of the column.

The default expression will be used in any insert operation that does not specify a value for the column. If there is no default for a column, then the default is null.

UNIQUE (column constraint)

UNIQUE ( column_name [, ... ] ) (table constraint)

The UNIQUE constraint specifies that a group of one or more columns of a table can contain only unique values. The behavior of the unique table constraint is the same as that for column constraints, with the additional capability to span multiple columns.

For the purpose of a unique constraint, null values are not considered equal.

Each unique table constraint must name a set of columns that is different from the set of columns named by any other unique or primary key constraint defined for the table. (Otherwise it would just be the same constraint listed twice.)

PRIMARY KEY (column constraint)

PRIMARY KEY ( column_name [, ... ] ) (table constraint)

The PRIMARY KEY constraint specifies that a column or columns of a table can contain only unique (non-duplicate), nonnull values. Only one primary key can be specified for a table, whether as a column constraint or a table constraint.

The primary key constraint should name a set of columns that is different from the set of columns named by any unique constraint defined for the same table. (Otherwise, the unique constraint is redundant and will be discarded.)
**CREATE TABLE**

*PRIMARY KEY* enforces the same data constraints as a combination of *UNIQUE* and *NOT NULL*, but identifying a set of columns as the primary key also provides metadata about the design of the schema, since a primary key implies that other tables can rely on this set of columns as a unique identifier for rows.

**EXCLUDE [ USING index_method ] ( exclude_element WITH operator [, ... ] )**

*index_parameters* [ WHERE ( predicate ) ]

The *EXCLUDE* clause defines an exclusion constraint, which guarantees that if any two rows are compared on the specified column(s) or expression(s) using the specified operator(s), not all of these comparisons will return *TRUE*. If all of the specified operators test for equality, this is equivalent to a *UNIQUE* constraint, although an ordinary unique constraint will be faster. However, exclusion constraints can specify constraints that are more general than simple equality. For example, you can specify a constraint that no two rows in the table contain overlapping circles (see Section 8.8) by using the && operator.

Exclusion constraints are implemented using an index, so each specified operator must be associated with an appropriate operator class (see Section 11.9) for the index access method *index_method*. The operators are required to be commutative. Each *exclude_element* can optionally specify an operator class and/or ordering options; these are described fully under *CREATE INDEX*.

The *predicate* allows you to specify an exclusion constraint on a subset of the table; internally this creates a partial index. Note that parentheses are required around the predicate.

**REFERENCES reftable [ ( refcolumn ) ] [ MATCH matchtype ] [ ON DELETE action ] [ ON UPDATE action ] (column constraint)**

**FOREIGN KEY ( column_name [, ... ] ) REFERENCES reftable [ ( refcolumn [, ... ] ) ] [ MATCH matchtype ] [ ON DELETE action ] [ ON UPDATE action ] (table constraint)**

These clauses specify a foreign key constraint, which requires that a group of one or more columns of the new table must only contain values that match values in the referenced column(s) of some row of the referenced table. If the *refcolumn* list is omitted, the primary key of the *reftable* is used. The referenced columns must be the columns of a non-deferrable unique or primary key constraint in the referenced table. Note that foreign key constraints cannot be defined between temporary tables and permanent tables.

A value inserted into the referencing column(s) is matched against the values of the referenced table and referenced columns using the given match type. There are three match types: *MATCH FULL*, *MATCH PARTIAL*, and *MATCH SIMPLE* (which is the default). *MATCH FULL* will not allow one column of a multicoloumn foreign key to be null unless all foreign key columns are null; if they are all null, the row is not required to have a match in the referenced table. *MATCH SIMPLE* allows any of the foreign key columns to be null; if any of them are null, the row is not required to have a match in the referenced table. *MATCH PARTIAL* is not yet implemented. (Of course, *NOT NULL* constraints can be applied to the referencing column(s) to prevent these cases from arising.)

In addition, when the data in the referenced columns is changed, certain actions are performed on the data in this table’s columns. The *ON DELETE* clause specifies the action to perform when a referenced row in the referenced table is being deleted. Likewise, the *ON UPDATE* clause specifies the action to perform when a referenced column in the referenced table is being updated to a new...
value. If the row is updated, but the referenced column is not actually changed, no action is done.
Referential actions other than the NO ACTION check cannot be deferred, even if the constraint is declared deferrable. There are the following possible actions for each clause:

NO ACTION

Produce an error indicating that the deletion or update would create a foreign key constraint violation. If the constraint is deferred, this error will be produced at constraint check time if there still exist any referencing rows. This is the default action.

RESTRICT

Produce an error indicating that the deletion or update would create a foreign key constraint violation. This is the same as NO ACTION except that the check is not deferrable.

CASCADE

Delete any rows referencing the deleted row, or update the values of the referencing column(s) to the new values of the referenced columns, respectively.

SET NULL

Set the referencing column(s) to null.

SET DEFAULT

Set the referencing column(s) to their default values. (There must be a row in the referenced table matching the default values, if they are not null, or the operation will fail.)

If the referenced column(s) are changed frequently, it might be wise to add an index to the referencing column(s) so that referential actions associated with the foreign key constraint can be performed more efficiently.

DEFERRABLE

This controls whether the constraint can be deferred. A constraint that is not deferrable will be checked immediately after every command. Checking of constraints that are deferrable can be postponed until the end of the transaction (using the SET CONSTRAINTS command). NOT DEFERRABLE is the default. Currently, only UNIQUE, PRIMARY KEY, EXCLUDE, and REFERENCES (foreign key) constraints accept this clause. NOT NULL and CHECK constraints are not deferrable. Note that deferrable constraints cannot be used as conflict arbitrators in an INSERT statement that includes an ON CONFLICT DO UPDATE clause.

INITIALLY IMMEDIATE

INITIALLY DEFERRED

If a constraint is deferrable, this clause specifies the default time to check the constraint. If the constraint is INITIALLY IMMEDIATE, it is checked after each statement. This is the default. If the constraint is INITIALLY DEFERRED, it is checked only at the end of the transaction. The constraint check time can be altered with the SET CONSTRAINTS command.

WITH ( storage_parameter [= value] [, ... ] )

This clause specifies optional storage parameters for a table or index; see Storage Parameters for more information. The WITH clause for a table can also include OIDS=TRUE (or just OIDS) to specify that rows of the new table should have OIDs (object identifiers) assigned to them, or OIDS=FALSE to specify that the rows should not have OIDs. If OIDS is not specified, the default setting depends upon the default_with_oids configuration parameter. (If the new table
inherits from any tables that have OIDs, then OIDS=TRUE is forced even if the command says OIDS=FALSE.)

If OIDS=FALSE is specified or implied, the new table does not store OIDs and no OID will be assigned for a row inserted into it. This is generally considered worthwhile, since it will reduce OID consumption and thereby postpone the wraparound of the 32-bit OID counter. Once the counter wraps around, OIDs can no longer be assumed to be unique, which makes them considerably less useful. In addition, excluding OIDs from a table reduces the space required to store the table on disk by 4 bytes per row (on most machines), slightly improving performance.

To remove OIDs from a table after it has been created, use ALTER TABLE.

WITH OIDS
WITHOUT OIDS

These are obsolescent syntaxes equivalent to WITH (OIDS) and WITH (OIDS=FALSE), respectively. If you wish to give both an OIDS setting and storage parameters, you must use the WITH (...) syntax; see above.

ON COMMIT

The behavior of temporary tables at the end of a transaction block can be controlled using ON COMMIT. The three options are:

PRESERVE ROWS

No special action is taken at the ends of transactions. This is the default behavior.

DELETE ROWS

All rows in the temporary table will be deleted at the end of each transaction block. Essentially, an automatic TRUNCATE is done at each commit.

DROP

The temporary table will be dropped at the end of the current transaction block.

TABLESPACE tablespace_name

The tablespace_name is the name of the tablespace in which the new table is to be created. If not specified, default_tablespace is consulted, or temp_tablespaces if the table is temporary.

USING INDEX TABLESPACE tablespace_name

This clause allows selection of the tablespace in which the index associated with a UNIQUE, PRIMARY KEY, or EXCLUDE constraint will be created. If not specified, default_tablespace is consulted, or temp_tablespaces if the table is temporary.

Storage Parameters

The WITH clause can specify storage parameters for tables, and for indexes associated with a UNIQUE, PRIMARY KEY, or EXCLUDE constraint. Storage parameters for indexes are documented in CREATE INDEX. The storage parameters currently available for tables are listed below. For many of these parameters, as shown, there is an additional parameter with the same name prefixed with toast., which controls the behavior of the table’s secondary TOAST table, if any (see Section 65.2 for more information about TOAST). If a table parameter value is set and the equivalent toast. parameter is not, the TOAST table will use the table’s parameter value.
CREATE TABLE

fillfactor (integer)

The fillfactor for a table is a percentage between 10 and 100. 100 (complete packing) is the default. When a smaller fillfactor is specified, INSERT operations pack table pages only to the indicated percentage; the remaining space on each page is reserved for updating rows on that page. This gives UPDATE a chance to place the updated copy of a row on the same page as the original, which is more efficient than placing it on a different page. For a table whose entries are never updated, complete packing is the best choice, but in heavily updated tables smaller fillfactors are appropriate. This parameter cannot be set for TOAST tables.

parallel_workers (integer)

This sets the number of workers that should be used to assist a parallel scan of this table. If not set, the system will determine a value based on the relation size. The actual number of workers chosen by the planner may be less, for example due to the setting of max_worker_processes.

autovacuum_enabled, toast.autovacuum_enabled (boolean)

Enables or disables the autovacuum daemon for a particular table. If true, the autovacuum daemon will perform automatic VACUUM and/or ANALYZE operations on this table following the rules discussed in Section 24.1.6. If false, this table will not be autovacuumed, except to prevent transaction ID wraparound. See Section 24.1.5 for more about wraparound prevention. Note that the autovacuum daemon does not run at all (except to prevent transaction ID wraparound) if the autovacuum parameter is false; setting individual tables’ storage parameters does not override that. Therefore there is seldom much point in explicitly setting this storage parameter to true, only to false.

autovacuum_vacuum_threshold, toast.autovacuum_vacuum_threshold (integer)

Per-table value for autovacuum_vacuum_threshold parameter.

autovacuum_vacuum_scale_factor, toast.autovacuum_vacuum_scale_factor (float4)

Per-table value for autovacuum_vacuum_scale_factor parameter.

autovacuum_analyze_threshold, toast.autovacuum_analyze_threshold (integer)

Per-table value for autovacuum_analyze_threshold parameter.

autovacuum_analyze_scale_factor (float4)

Per-table value for autovacuum_analyze_scale_factor parameter.

autovacuum_vacuum_cost_delay, toast.autovacuum_vacuum_cost_delay (integer)

Per-table value for autovacuum_vacuum_cost_delay parameter.

autovacuum_vacuum_cost_limit, toast.autovacuum_vacuum_cost_limit (integer)

Per-table value for autovacuum_vacuum_cost_limit parameter.

autovacuum_freeze_min_age, toast.autovacuum_freeze_min_age (integer)

Per-table value for vacuum_freeze_min_age parameter. Note that autovacuum will ignore per-table autovacuum_freeze_min_age parameters that are larger than half the system-wide autovacuum_freeze_max_age setting.

autovacuum_freeze_max_age, toast.autovacuum_freeze_max_age (integer)

Per-table value for autovacuum_freeze_max_age parameter. Note that autovacuum will ignore per-table autovacuum_freeze_max_age parameters that are larger than the system-wide setting (it can only be set smaller).

autovacuum_freeze_table_age, toast.autovacuum_freeze_table_age (integer)

Per-table value for vacuum_freeze_table_age parameter.
CREATE TABLE

autovacuum_multixact_freeze_min_age, toast.autovacuum_multixact_freeze_min_age (integer)

Per-table value for vacuum_multixact_freeze_min_age parameter. Note that autovacuum will ignore per-table autovacuum_multixact_freeze_min_age parameters that are larger than half the system-wide autovacuum_multixact_freeze_max_age setting.

autovacuum_multixact_freeze_max_age, toast.autovacuum_multixact_freeze_max_age (integer)

Per-table value for autovacuum_multixact_freeze_max_age parameter. Note that autovacuum will ignore per-table autovacuum_multixact_freeze_max_age parameters that are larger than the system-wide setting (it can only be set smaller).

autovacuum_multixact_freeze_table_age, toast.autovacuum_multixact_freeze_table_age (integer)

Per-table value for vacuum_multixact_freeze_table_age parameter.

log_autovacuum_min_duration, toast.log_autovacuum_min_duration (integer)

Per-table value for log_autovacuum_min_duration parameter.

user_catalog_table (boolean)

Declare the table as an additional catalog table for purposes of logical replication. See Section 47.6.2 for details. This parameter cannot be set for TOAST tables.

Notes

Using OIDs in new applications is not recommended: where possible, using a SERIAL or other sequence generator as the table’s primary key is preferred. However, if your application does make use of OIDs to identify specific rows of a table, it is recommended to create a unique constraint on the oid column of that table, to ensure that OIDs in the table will indeed uniquely identify rows even after counter wraparound. Avoid assuming that OIDs are unique across tables; if you need a database-wide unique identifier, use the combination of tableoid and row OID for the purpose.

Tip: The use of OIDS=FALSE is not recommended for tables with no primary key, since without either an OID or a unique data key, it is difficult to identify specific rows.

PostgreSQL automatically creates an index for each unique constraint and primary key constraint to enforce uniqueness. Thus, it is not necessary to create an index explicitly for primary key columns. (See CREATE INDEX for more information.)

Unique constraints and primary keys are not inherited in the current implementation. This makes the combination of inheritance and unique constraints rather dysfunctional.

A table cannot have more than 1600 columns. (In practice, the effective limit is usually lower because of tuple-length constraints.)

Examples

Create table films and table distributors:

CREATE TABLE films (
CREATE TABLE

code char(5) CONSTRAINT firstkey PRIMARY KEY,
title varchar(40) NOT NULL,
did integer NOT NULL,
date_prod date,
kind varchar(10),
len interval hour to minute
);

CREATE TABLE distributors {
    did integer PRIMARY KEY DEFAULT nextval('serial'),
    name varchar(40) NOT NULL CHECK (name <> "")
};

Create a table with a 2-dimensional array:

CREATE TABLE array_int (
    vector int[][]
);
CREATE TABLE code char(5),
title varchar(40),
did integer,
date_prod date,
kind varchar(10),
len interval hour to minute,
CONSTRAINT code_title PRIMARY KEY(code,title)
);

Define a primary key constraint for table distributors. The following two examples are equivalent, the first using the table constraint syntax, the second the column constraint syntax:

CREATE TABLE distributors (  
did integer,
name varchar(40),
PRIMARY KEY(did)
);

CREATE TABLE distributors (  
did integer PRIMARY KEY,
name varchar(40)
);

Assign a literal constant default value for the column name, arrange for the default value of column did to be generated by selecting the next value of a sequence object, and make the default value of modtime be the time at which the row is inserted:

CREATE TABLE distributors (  
name varchar(40) DEFAULT 'Luso Films',
did integer DEFAULT nextval('distributors_serial'),
modtime timestamp DEFAULT current_timestamp
);

Define two NOT NULL column constraints on the table distributors, one of which is explicitly given a name:

CREATE TABLE distributors (  
did integer CONSTRAINT no_null NOT NULL,
name varchar(40) NOT NULL
);

Define a unique constraint for the name column:

CREATE TABLE distributors (  
did integer,
name varchar(40) UNIQUE
);

The same, specified as a table constraint:

CREATE TABLE distributors (  
did integer,
CREATE TABLE
    name varchar(40),
    UNIQUE(name)
);

Create the same table, specifying 70% fill factor for both the table and its unique index:

CREATE TABLE distributors (
    did integer,
    name varchar(40),
    UNIQUE(name) WITH (fillfactor=70)
)
WITH (fillfactor=70);

Create table circles with an exclusion constraint that prevents any two circles from overlapping:

CREATE TABLE circles (
    c circle,
    EXCLUDE USING gist (c WITH &&)
);

Create table cinemas in tablespace diskvol1:

CREATE TABLE cinemas (
    id serial,
    name text,
    location text
) TABLESPACE diskvol1;

Create a composite type and a typed table:

CREATE TYPE employee_type AS (name text, salary numeric);

CREATE TABLE employees OF employee_type (
    PRIMARY KEY (name),
    salary WITH OPTIONS DEFAULT 1000
);

Compatibility

The CREATE TABLE command conforms to the SQL standard, with exceptions listed below.

Temporary Tables

Although the syntax of CREATE TEMPORARY TABLE resembles that of the SQL standard, the effect is not the same. In the standard, temporary tables are defined just once and automatically exist (starting with empty contents) in every session that needs them. PostgreSQL instead requires each session to issue its own CREATE TEMPORARY TABLE command for each temporary table to be used. This allows
different sessions to use the same temporary table name for different purposes, whereas the standard’s
approach constrains all instances of a given temporary table name to have the same table structure.
The standard’s definition of the behavior of temporary tables is widely ignored. PostgreSQL’s behav-
ior on this point is similar to that of several other SQL databases.
The SQL standard also distinguishes between global and local temporary tables, where a local tempo-
rary table has a separate set of contents for each SQL module within each session, though its definition
is still shared across sessions. Since PostgreSQL does not support SQL modules, this distinction is
not relevant in PostgreSQL.
For compatibility’s sake, PostgreSQL will accept the GLOBAL and LOCAL keywords in a temporary
table declaration, but they currently have no effect. Use of these keywords is discouraged, since future
versions of PostgreSQL might adopt a more standard-compliant interpretation of their meaning.
The ON COMMIT clause for temporary tables also resembles the SQL standard, but has some differ-
ences. If the ON COMMIT clause is omitted, SQL specifies that the default behavior is ON COMMIT
DELETE ROWS. However, the default behavior in PostgreSQL is ON COMMIT PRESERVE ROWS. The
ON COMMIT DROP option does not exist in SQL.

Non-deferred Uniqueness Constraints
When a UNIQUE or PRIMARY KEY constraint is not deferrable, PostgreSQL checks for uniqueness
immediately whenever a row is inserted or modified. The SQL standard says that uniqueness should
be enforced only at the end of the statement; this makes a difference when, for example, a single com-
mmand updates multiple key values. To obtain standard-compliant behavior, declare the constraint as
DEFERRABLE but not deferred (i.e., INITIALLY IMMEDIATE). Be aware that this can be significantly
slower than immediate uniqueness checking.

Column Check Constraints
The SQL standard says that CHECK column constraints can only refer to the column they apply to; only
CHECK table constraints can refer to multiple columns. PostgreSQL does not enforce this restriction;
it treats column and table check constraints alike.

EXCLUDE Constraint
The EXCLUDE constraint type is a PostgreSQL extension.

NULL “Constraint”
The NULL “constraint” (actually a non-constraint) is a PostgreSQL extension to the SQL standard
that is included for compatibility with some other database systems (and for symmetry with the NOT
NULL constraint). Since it is the default for any column, its presence is simply noise.

Inheritance
Multiple inheritance via the INHERITS clause is a PostgreSQL language extension. SQL:1999 and
later define single inheritance using a different syntax and different semantics. SQL:1999-style inher-
itance is not yet supported by PostgreSQL.
Zero-column Tables

PostgreSQL allows a table of no columns to be created (for example, `CREATE TABLE foo();`). This is an extension from the SQL standard, which does not allow zero-column tables. Zero-column tables are not in themselves very useful, but disallowing them creates odd special cases for `ALTER TABLE DROP COLUMN`, so it seems cleaner to ignore this spec restriction.

LIKE Clause

While a LIKE clause exists in the SQL standard, many of the options that PostgreSQL accepts for it are not in the standard, and some of the standard’s options are not implemented by PostgreSQL.

WITH Clause

The WITH clause is a PostgreSQL extension; neither storage parameters nor OIDs are in the standard.

Tablespaces

The PostgreSQL concept of tablespaces is not part of the standard. Hence, the clauses `TABLESPACE` and `USING INDEX TABLESPACE` are extensions.

Typed Tables

Typed tables implement a subset of the SQL standard. According to the standard, a typed table has columns corresponding to the underlying composite type as well as one other column that is the “self-referencing column”. PostgreSQL does not support these self-referencing columns explicitly, but the same effect can be had using the OID feature.

See Also

`ALTER TABLE`, `DROP TABLE`, `CREATE TABLE AS`, `CREATE TABLESPACE`, `CREATE TYPE`
CREATE TABLE AS

Name
CREATE TABLE AS — define a new table from the results of a query

Synopsis
CREATE [ [ GLOBAL | LOCAL ] { TEMPORARY | TEMP } | UNLOGGED ] TABLE [ IF NOT EXISTS ] table_name
{ (column_name [, ...] ) }
[ WITH ( storage_parameter [= value] [, ... ] ) | WITH OIDS | WITHOUT OIDS ]
[ ON COMMIT { PRESERVE ROWS | DELETE ROWS | DROP } ]
[ TABLESPACE tablespace_name ]
AS query
[ WITH [ NO ] DATA ]

Description
CREATE TABLE AS creates a table and fills it with data computed by a SELECT command. The table
columns have the names and data types associated with the output columns of the SELECT (except
that you can override the column names by giving an explicit list of new column names).

CREATE TABLE AS bears some resemblance to creating a view, but it is really quite different: it
creates a new table and evaluates the query just once to fill the new table initially. The new table will
not track subsequent changes to the source tables of the query. In contrast, a view re-evaluates its
defining SELECT statement whenever it is queried.

Parameters
GLOBAL or LOCAL
Ignored for compatibility. Use of these keywords is deprecated; refer to CREATE TABLE for
details.

TEMPORARY or TEMP
If specified, the table is created as a temporary table. Refer to CREATE TABLE for details.

UNLOGGED
If specified, the table is created as an unlogged table. Refer to CREATE TABLE for details.

IF NOT EXISTS
Do not throw an error if a relation with the same name already exists. A notice is issued in this
case. Refer to CREATE TABLE for details.

table_name
The name (optionally schema-qualified) of the table to be created.

column_name
The name of a column in the new table. If column names are not provided, they are taken from
the output column names of the query.
CREATE TABLE AS

WITH ( storage_parameter [= value] [, ... ] )

This clause specifies optional storage parameters for the new table; see Storage Parameters for more information. The WITH clause can also include \texttt{OIDS=TRUE} (or just \texttt{OIDS}) to specify that rows of the new table should have OIDs (object identifiers) assigned to them, or \texttt{OIDS=FALSE} to specify that the rows should not have OIDs. See CREATE TABLE for more information.

\textbf{WITH OIDS}

\textbf{WITHOUT OIDS}

These are obsolescent syntaxes equivalent to \texttt{WITH (OIDS)} and \texttt{WITH (OIDS=FALSE)}, respectively. If you wish to give both an OIDs setting and storage parameters, you must use the \texttt{WITH ( ... )} syntax; see above.

\textbf{ON COMMIT}

The behavior of temporary tables at the end of a transaction block can be controlled using \texttt{ON COMMIT}. The three options are:

\textbf{PRESERVE ROWS}

No special action is taken at the ends of transactions. This is the default behavior.

\textbf{DELETE ROWS}

All rows in the temporary table will be deleted at the end of each transaction block. Essentially, an automatic TRUNCATE is done at each commit.

\textbf{DROP}

The temporary table will be dropped at the end of the current transaction block.

\textbf{TABLESPACE} tablespace_name

The \texttt{tablespace_name} is the name of the tablespace in which the new table is to be created. If not specified, \texttt{default_tablespace} is consulted, or \texttt{temp_tablespaces} if the table is temporary.

\textbf{query}

A SELECT, TABLE, or VALUES command, or an EXECUTE command that runs a prepared SELECT, TABLE, or VALUES query.

\textbf{WITH [ NO ] DATA}

This clause specifies whether or not the data produced by the query should be copied into the new table. If not, only the table structure is copied. The default is to copy the data.

\textbf{Notes}

This command is functionally similar to SELECT INTO, but it is preferred since it is less likely to be confused with other uses of the \texttt{SELECT INTO} syntax. Furthermore, CREATE TABLE AS offers a superset of the functionality offered by SELECT INTO.

The CREATE TABLE AS command allows the user to explicitly specify whether OIDs should be included. If the presence of OIDs is not explicitly specified, the default_with_Oids configuration variable is used.
Examples

Create a new table `films_recent` consisting of only recent entries from the table `films`:

```sql
CREATE TABLE films_recent AS
    SELECT * FROM films WHERE date_prod >= '2002-01-01';
```

To copy a table completely, the short form using the `TABLE` command can also be used:

```sql
CREATE TABLE films2 AS
    TABLE films;
```

Create a new temporary table `films_recent`, consisting of only recent entries from the table `films`, using a prepared statement. The new table has OIDs and will be dropped at commit:

```sql
PREPARE recentfilms(date) AS
    SELECT * FROM films WHERE date_prod > $1;
CREATE TEMP TABLE films_recent WITH (OIDS) ON COMMIT DROP AS
    EXECUTE recentfilms('2002-01-01');
```

Compatibility

`CREATE TABLE AS` conforms to the SQL standard. The following are nonstandard extensions:

- The standard requires parentheses around the subquery clause; in PostgreSQL, these parentheses are optional.
- In the standard, the `WITH [ NO ] DATA` clause is required; in PostgreSQL it is optional.
- PostgreSQL handles temporary tables in a way rather different from the standard; see `CREATE TABLE` for details.
- The `WITH` clause is a PostgreSQL extension; neither storage parameters nor OIDs are in the standard.
- The PostgreSQL concept of tablespaces is not part of the standard. Hence, the clause `TABLESPACE` is an extension.

See Also

`CREATE MATERIALIZED VIEW, CREATE TABLE, EXECUTE, SELECT, SELECT INTO, VALUES`
CREATE TABLESPACE

Name
CREATE TABLESPACE — define a new tablespace

Synopsis
CREATE TABLESPACE tablespace_name
    [ OWNER { new_owner | CURRENT_USER | SESSION_USER } ]
    LOCATION 'directory'
    [ WITH ( tablespace_option = value [, ... ] ) ]

Description
CREATE TABLESPACE registers a new cluster-wide tablespace. The tablespace name must be distinct from the name of any existing tablespace in the database cluster.

A tablespace allows superusers to define an alternative location on the file system where the data files containing database objects (such as tables and indexes) can reside.

A user with appropriate privileges can pass tablespace_name to CREATE DATABASE, CREATE TABLE, CREATE INDEX or ADD CONSTRAINT to have the data files for these objects stored within the specified tablespace.

Warning
A tablespace cannot be used independently of the cluster in which it is defined; see Section 22.6.

Parameters

tablespace_name
The name of a tablespace to be created. The name cannot begin with pg_, as such names are reserved for system tablespaces.

user_name
The name of the user who will own the tablespace. If omitted, defaults to the user executing the command. Only superusers can create tablespaces, but they can assign ownership of tablespaces to non-superusers.

directory
The directory that will be used for the tables. The directory should be empty and must be owned by the PostgreSQL system user. The directory must be specified by an absolute path name.
tablespace_option

A tablespace parameter to be set or reset. Currently, the only available parameters are seq_page_cost, random_page_cost and effective_io_concurrency. Setting either value for a particular tablespace will override the planner’s usual estimate of the cost of reading pages from tables in that tablespace, as established by the configuration parameters of the same name (see seq_page_cost, random_page_cost, effective_io_concurrency). This may be useful if one tablespace is located on a disk which is faster or slower than the remainder of the I/O subsystem.

Notes

Tablespaces are only supported on systems that support symbolic links. CREATE TABLESPACE cannot be executed inside a transaction block.

Examples

Create a tablespace dbspace at /data/dbs:

CREATE TABLESPACE dbspace LOCATION '/data/dbs';

Create a tablespace indexspace at /data/indexes owned by user genevieve:

CREATE TABLESPACE indexspace OWNER genevieve LOCATION '/data/indexes';

Compatibility

CREATE TABLESPACE is a PostgreSQL extension.

See Also

CREATE DATABASE, CREATE TABLE, CREATE INDEX, DROP TABLESPACE, ALTER TABLESPACE
CREATE TEXT SEARCH CONFIGURATION

Name
CREATE TEXT SEARCH CONFIGURATION — define a new text search configuration

Synopsis
CREATE TEXT SEARCH CONFIGURATION name ( 
    PARSER = parser_name | 
    COPY = source_config 
)

Description
CREATE TEXT SEARCH CONFIGURATION creates a new text search configuration. A text search configuration specifies a text search parser that can divide a string into tokens, plus dictionaries that can be used to determine which tokens are of interest for searching.

If only the parser is specified, then the new text search configuration initially has no mappings from token types to dictionaries, and therefore will ignore all words. Subsequent ALTER TEXT SEARCH CONFIGURATION commands must be used to create mappings to make the configuration useful. Alternatively, an existing text search configuration can be copied.

If a schema name is given then the text search configuration is created in the specified schema. Otherwise it is created in the current schema.

The user who defines a text search configuration becomes its owner.

Refer to Chapter 12 for further information.

Parameters

name

The name of the text search configuration to be created. The name can be schema-qualified.

parser_name

The name of the text search parser to use for this configuration.

source_config

The name of an existing text search configuration to copy.

Notes

The PARSER and COPY options are mutually exclusive, because when an existing configuration is copied, its parser selection is copied too.
CREATE TEXT SEARCH CONFIGURATION

Compatibility

There is no CREATE TEXT SEARCH CONFIGURATION statement in the SQL standard.

See Also

ALTER TEXT SEARCH CONFIGURATION, DROP TEXT SEARCH CONFIGURATION
CREATE TEXT SEARCH DICTIONARY

Name
CREATE TEXT SEARCH DICTIONARY — define a new text search dictionary

Synopsis
CREATE TEXT SEARCH DICTIONARY name ( 
  TEMPLATE = template 
  [, option = value [, ... ]] 
)

Description
CREATE TEXT SEARCH DICTIONARY creates a new text search dictionary. A text search dictionary specifies a way of recognizing interesting or uninteresting words for searching. A dictionary depends on a text search template, which specifies the functions that actually perform the work. Typically the dictionary provides some options that control the detailed behavior of the template’s functions.
If a schema name is given then the text search dictionary is created in the specified schema. Otherwise it is created in the current schema.
The user who defines a text search dictionary becomes its owner.
Refer to Chapter 12 for further information.

Parameters
name
  The name of the text search dictionary to be created. The name can be schema-qualified.
template
  The name of the text search template that will define the basic behavior of this dictionary.
option
  The name of a template-specific option to be set for this dictionary.
value
  The value to use for a template-specific option. If the value is not a simple identifier or number, it must be quoted (but you can always quote it, if you wish).
The options can appear in any order.

Examples
The following example command creates a Snowball-based dictionary with a nonstandard list of stop words.
CREATE TEXT SEARCH DICTIONARY my_russian (
CREATE TEXT SEARCH DICTIONARY

template = snowball,
language = russian,
stopwords = myrussian
);

Compatibility

There is no CREATE TEXT SEARCH DICTIONARY statement in the SQL standard.

See Also

ALTER TEXT SEARCH DICTIONARY, DROP TEXT SEARCH DICTIONARY
CREATE TEXT SEARCH PARSER

Name

CREATE TEXT SEARCH PARSER — define a new text search parser

Synopsis

CREATE TEXT SEARCH PARSER name (  
    START = start_function ,  
    GETTOKEN = gettoken_function ,  
    END = end_function ,  
    LEXTYPES = lextypes_function  
    [, HEADLINE = headline_function ]
)

Description

CREATE TEXT SEARCH PARSER creates a new text search parser. A text search parser defines a  
method for splitting a text string into tokens and assigning types (categories) to the tokens. A parser  
is not particularly useful by itself, but must be bound into a text search configuration along with some  
text search dictionaries to be used for searching.  

If a schema name is given then the text search parser is created in the specified schema. Otherwise it  
is created in the current schema.  

You must be a superuser to use CREATE TEXT SEARCH PARSER. (This restriction is made because  
an erroneous text search parser definition could confuse or even crash the server.)  
Refer to Chapter 12 for further information.

Parameters

name

The name of the text search parser to be created. The name can be schema-qualified.

start_function

The name of the start function for the parser.

gettoken_function

The name of the get-next-token function for the parser.

end_function

The name of the end function for the parser.

lextypes_function

The name of the lextypes function for the parser (a function that returns information about the  
set of token types it produces).
CREATE TEXT SEARCH PARSER

headline_function

The name of the headline function for the parser (a function that summarizes a set of tokens). The function names can be schema-qualified if necessary. Argument types are not given, since the argument list for each type of function is predetermined. All except the headline function are required. The arguments can appear in any order, not only the one shown above.

Compatibility
There is no CREATE TEXT SEARCH PARSER statement in the SQL standard.

See Also
ALTER TEXT SEARCH PARSER, DROP TEXT SEARCH PARSER
CREATE TEXT SEARCH TEMPLATE

Name
CREATE TEXT SEARCH TEMPLATE — define a new text search template

Synopsis
CREATE TEXT SEARCH TEMPLATE name (  
  [ INIT = init_function , ]  
  LEXIZE = lexize_function  
)

Description
CREATE TEXT SEARCH TEMPLATE creates a new text search template. Text search templates define  
the functions that implement text search dictionaries. A template is not useful by itself, but must be  
instantiated as a dictionary to be used. The dictionary typically specifies parameters to be given to the  
template functions.
If a schema name is given then the text search template is created in the specified schema. Otherwise  
it is created in the current schema.
You must be a superuser to use CREATE TEXT SEARCH TEMPLATE. This restriction is made because  
an erroneous text search template definition could confuse or even crash the server. The reason for  
separating templates from dictionaries is that a template encapsulates the “unsafe” aspects of defining  
a dictionary. The parameters that can be set when defining a dictionary are safe for unprivileged users  
to set, and so creating a dictionary need not be a privileged operation.
Refer to Chapter 12 for further information.

Parameters

name
The name of the text search template to be created. The name can be schema-qualified.

init_function
The name of the init function for the template.

lexize_function
The name of the lexize function for the template.
The function names can be schema-qualified if necessary. Argument types are not given, since the  
argument list for each type of function is predetermined. The lexize function is required, but the init  
function is optional.
The arguments can appear in any order, not only the one shown above.
CREATE TEXT SEARCH TEMPLATE

Compatibility
There is no CREATE TEXT SEARCH TEMPLATE statement in the SQL standard.

See Also
ALTER TEXT SEARCH TEMPLATE, DROP TEXT SEARCH TEMPLATE
CREATE TRANSFORM

Name

CREATE TRANSFORM — define a new transform

Synopsis

CREATE [ OR REPLACE ] TRANSFORM FOR type_name LANGUAGE lang_name (  
    FROM SQL WITH FUNCTION from_sql_function_name (argument_type [, ...]),  
    TO SQL WITH FUNCTION to_sql_function_name (argument_type [, ...])  
);  

Description

CREATE TRANSFORM defines a new transform. CREATE OR REPLACE TRANSFORM will either create a new transform, or replace an existing definition.

A transform specifies how to adapt a data type to a procedural language. For example, when writing a function in PL/Python using the hstore type, PL/Python has no prior knowledge how to present hstore values in the Python environment. Language implementations usually default to using the text representation, but that is inconvenient when, for example, an associative array or a list would be more appropriate.

A transform specifies two functions:

• A “from SQL” function that converts the type from the SQL environment to the language. This function will be invoked on the arguments of a function written in the language.

• A “to SQL” function that converts the type from the language to the SQL environment. This function will be invoked on the return value of a function written in the language.

It is not necessary to provide both of these functions. If one is not specified, the language-specific default behavior will be used if necessary. (To prevent a transformation in a certain direction from happening at all, you could also write a transform function that always errors out.)

To be able to create a transform, you must own and have USAGE privilege on the type, have USAGE privilege on the language, and own and have EXECUTE privilege on the from-SQL and to-SQL functions, if specified.

Parameters

type_name

The name of the data type of the transform.

lang_name

The name of the language of the transform.
CREATE TRANSFORM

from_sql_function_name(argument_type [, ...])

The name of the function for converting the type from the SQL environment to the language. It must take one argument of type internal and return type internal. The actual argument will be of the type for the transform, and the function should be coded as if it were. (But it is not allowed to declare an SQL-level function returning internal without at least one argument of type internal.) The actual return value will be something specific to the language implementation.

to_sql_function_name(argument_type [, ...])

The name of the function for converting the type from the language to the SQL environment. It must take one argument of type internal and return the type that is the type for the transform. The actual argument value will be something specific to the language implementation.

Notes

Use DROP TRANSFORM to remove transforms.

Examples

To create a transform for type hstore and language plpythonu, first set up the type and the language:

CREATE TYPE hstore ...;

CREATE LANGUAGE plpythonu ...;

Then create the necessary functions:

CREATE FUNCTION hstore_to_plpython(val internal) RETURNS internal LANGUAGE C STRICT IMMUTABLE AS ...;

CREATE FUNCTION plpython_to_hstore(val internal) RETURNS hstore LANGUAGE C STRICT IMMUTABLE AS ...;

And finally create the transform to connect them all together:

CREATE TRANSFORM FOR hstore LANGUAGE plpythonu (  FROM SQL WITH FUNCTION hstore_to_plpython(internal),  TO SQL WITH FUNCTION plpython_to_hstore(internal) ) ;

In practice, these commands would be wrapped up in extensions.

The contrib section contains a number of extensions that provide transforms, which can serve as real-world examples.
Compatibility

This form of `CREATE TRANSFORM` is a PostgreSQL extension. There is a `CREATE TRANSFORM` command in the SQL standard, but it is for adapting data types to client languages. That usage is not supported by PostgreSQL.

See Also

`CREATE FUNCTION`, `CREATE LANGUAGE`, `CREATE TYPE`, `DROP TRANSFORM`
CREATE TRIGGER

Name
CREATE TRIGGER — define a new trigger

Synopsis
CREATE [ CONSTRAINT ] TRIGGER name { BEFORE | AFTER | INSTEAD OF } { event [ OR ... ] } 
ON table_name 
[ FROM referenced_table_name ]
[ NOT DEFERRABLE | [ DEFERRABLE ] [ INITIALLY IMMEDIATE | INITIALLY DEFERRED ] ]
[ FOR [ EACH ] { ROW | STATEMENT } ]
[ WHEN ( condition ) ]
EXECUTE PROCEDURE function_name ( arguments )

where event can be one of:

   INSERT
   UPDATE [ OF column_name [, ... ] ]
   DELETE
   TRUNCATE

Description
CREATE TRIGGER creates a new trigger. The trigger will be associated with the specified table, view, 
or foreign table and will execute the specified function function_name when certain events occur.

The trigger can be specified to fire before the operation is attempted on a row (before constraints are 
checked and the INSERT, UPDATE, or DELETE is attempted); or after the operation has completed 
(after constraints are checked and the INSERT, UPDATE, or DELETE has completed); or instead of the 
operation (in the case of inserts, updates or deletes on a view). If the trigger fires before or instead of 
the event, the trigger can skip the operation for the current row, or change the row being inserted (for 
INSERT and UPDATE operations only). If the trigger fires after the event, all changes, including the 
effects of other triggers, are “visible” to the trigger.

A trigger that is marked FOR EACH ROW is called once for every row that the operation modifies. For 
example, a DELETE that affects 10 rows will cause any ON DELETE triggers on the target relation to 
be called 10 separate times, once for each deleted row. In contrast, a trigger that is marked FOR EACH 
STATEMENT only executes once for any given operation, regardless of how many rows it modifies (in 
particular, an operation that modifies zero rows will still result in the execution of any applicable FOR 
EACH STATEMENT triggers). Note that with an INSERT with an ON CONFLICT DO UPDATE clause, 
both INSERT and UPDATE statement level trigger will be fired.

Triggers that are specified to fire INSTEAD OF the trigger event must be marked FOR EACH ROW, and 
can only be defined on views. BEFORE and AFTER triggers on a view must be marked as FOR EACH 
STATEMENT.

In addition, triggers may be defined to fire for TRUNCATE, though only FOR EACH STATEMENT.

The following table summarizes which types of triggers may be used on tables, views, and foreign tables:
CREATE TRIGGER

<table>
<thead>
<tr>
<th>When</th>
<th>Event</th>
<th>Row-level</th>
<th>Statement-level</th>
</tr>
</thead>
<tbody>
<tr>
<td>BEFORE</td>
<td>INSERT/UPDATE/DELETE</td>
<td>Tables and foreign tables</td>
<td>Tables, views, and foreign tables</td>
</tr>
<tr>
<td></td>
<td>TRUNCATE</td>
<td>—</td>
<td>Tables</td>
</tr>
<tr>
<td>AFTER</td>
<td>INSERT/UPDATE/DELETE</td>
<td>Tables and foreign tables</td>
<td>Tables, views, and foreign tables</td>
</tr>
<tr>
<td></td>
<td>TRUNCATE</td>
<td>—</td>
<td>Tables</td>
</tr>
<tr>
<td>INSTEAD OF</td>
<td>INSERT/UPDATE/DELETE</td>
<td>Views</td>
<td>—</td>
</tr>
<tr>
<td></td>
<td>TRUNCATE</td>
<td>—</td>
<td>—</td>
</tr>
</tbody>
</table>

Also, a trigger definition can specify a Boolean WHEN condition, which will be tested to see whether the trigger should be fired. In row-level triggers the WHEN condition can examine the old and/or new values of columns of the row. Statement-level triggers can also have WHEN conditions, although the feature is not so useful for them since the condition cannot refer to any values in the table.

If multiple triggers of the same kind are defined for the same event, they will be fired in alphabetical order by name.

When the CONSTRAINT option is specified, this command creates a constraint trigger. This is the same as a regular trigger except that the timing of the trigger firing can be adjusted using SET CONSTRAINTS. Constraint triggers must be AFTER ROW triggers on tables. They can be fired either at the end of the statement causing the triggering event, or at the end of the containing transaction; in the latter case they are said to be deferred. A pending deferred-trigger firing can also be forced to happen immediately by using SET CONSTRAINTS. Constraint triggers are expected to raise an exception when the constraints they implement are violated.

SELECT does not modify any rows so you cannot create SELECT triggers. Rules and views are more appropriate in such cases.

Refer to Chapter 37 for more information about triggers.

### Parameters

- **name**
  
  The name to give the new trigger. This must be distinct from the name of any other trigger for the same table. The name cannot be schema-qualified — the trigger inherits the schema of its table. For a constraint trigger, this is also the name to use when modifying the trigger’s behavior using SET CONSTRAINTS.

- **BEFORE**
- **AFTER**
- **INSTEAD OF**

  Determines whether the function is called before, after, or instead of the event. A constraint trigger can only be specified as AFTER.

- **event**

  One of INSERT, UPDATE, DELETE, or TRUNCATE; this specifies the event that will fire the trigger. Multiple events can be specified using OR.
CREATE TRIGGER

For UPDATE events, it is possible to specify a list of columns using this syntax:

```
UPDATE OF column_name1 [, column_name2 ... ]
```

The trigger will only fire if at least one of the listed columns is mentioned as a target of the UPDATE command.

INSTEAD OF UPDATE events do not support lists of columns.

table_name

The name (optionally schema-qualified) of the table, view, or foreign table the trigger is for.

referenced_table_name

The (possibly schema-qualified) name of another table referenced by the constraint. This option is used for foreign-key constraints and is not recommended for general use. This can only be specified for constraint triggers.

DEFERRABLE
NOT DEFERRABLE
INITIALLY IMMEDIATE
INITIALLY DEFERRED

The default timing of the trigger. See the CREATE TABLE documentation for details of these constraint options. This can only be specified for constraint triggers.

FOR EACH ROW
FOR EACH STATEMENT

This specifies whether the trigger procedure should be fired once for every row affected by the trigger event, or just once per SQL statement. If neither is specified, FOR EACH STATEMENT is the default. Constraint triggers can only be specified FOR EACH ROW.

condition

A Boolean expression that determines whether the trigger function will actually be executed. If WHEN is specified, the function will only be called if the condition returns true. In FOR EACH ROW triggers, the WHEN condition can refer to columns of the old and/or new row values by writing OLD.column_name or NEW.column_name respectively. Of course, INSERT triggers cannot refer to OLD and DELETE triggers cannot refer to NEW.

INSTEAD OF triggers do not support WHEN conditions.

Currently, WHEN expressions cannot contain subqueries.

Note that for constraint triggers, evaluation of the WHEN condition is not deferred, but occurs immediately after the row update operation is performed. If the condition does not evaluate to true then the trigger is not queued for deferred execution.

function_name

A user-supplied function that is declared as taking no arguments and returning type trigger, which is executed when the trigger fires.

arguments

An optional comma-separated list of arguments to be provided to the function when the trigger is executed. The arguments are literal string constants. Simple names and numeric constants can be written here, too, but they will all be converted to strings. Please check the description of the implementation language of the trigger function to find out how these arguments can be accessed within the function; it might be different from normal function arguments.
**CREATE TRIGGER**

To create a trigger on a table, the user must have the **TRIGGER** privilege on the table. The user must also have **EXECUTE** privilege on the trigger function.

Use **DROP TRIGGER** to remove a trigger.

A column-specific trigger (one defined using the **UPDATE** OF column_name syntax) will fire when any of its columns are listed as targets in the **UPDATE** command’s **SET** list. It is possible for a column’s value to change even when the trigger is not fired, because changes made to the row’s contents by **BEFORE UPDATE** triggers are not considered. Conversely, a command such as **UPDATE ... SET x = x ...** will fire a trigger on column x, even though the column’s value did not change.

In a **BEFORE** trigger, the **WHEN** condition is evaluated just before the function is or would be executed, so using **WHEN** is not materially different from testing the same condition at the beginning of the trigger function. Note in particular that the **NEW** row seen by the condition is the current value, as possibly modified by earlier triggers. Also, a **BEFORE** trigger’s **WHEN** condition is not allowed to examine the system columns of the **NEW** row (such as **oid**), because those won’t have been set yet.

In an **AFTER** trigger, the **WHEN** condition is evaluated just after the row update occurs, and it determines whether an event is queued to fire the trigger at the end of statement. So when an **AFTER** trigger’s **WHEN** condition does not return true, it is not necessary to queue an event nor to re-fetch the row at end of statement. This can result in significant speedups in statements that modify many rows, if the trigger only needs to be fired for a few of the rows.

Statement-level triggers on a view are fired only if the action on the view is handled by a row-level **INSTEAD OF** trigger. If the action is handled by an **INSTEAD** rule, then whatever statements are emitted by the rule are executed in place of the original statement naming the view, so that the triggers that will be fired are those on tables named in the replacement statements. Similarly, if the view is automatically updatable, then the action is handled by automatically rewriting the statement into an action on the view’s base table, so that the base table’s statement-level triggers are the ones that are fired.

In PostgreSQL versions before 7.3, it was necessary to declare trigger functions as returning the placeholder type **opaque**, rather than **trigger**. To support loading of old dump files, **CREATE TRIGGER** will accept a function declared as returning **opaque**, but it will issue a notice and change the function’s declared return type to **trigger**.

### Examples

**Execute the function** check_account_update **whenever a row of the table accounts is about to be updated:**

```
CREATE TRIGGER check_update
    BEFORE UPDATE ON accounts
    FOR EACH ROW
    EXECUTE PROCEDURE check_account_update();
```

**The same, but only execute the function if column balance is specified as a target in the UPDATE command:**

```
CREATE TRIGGER check_update
    BEFORE UPDATE OF balance ON accounts
    FOR EACH ROW
    EXECUTE PROCEDURE check_account_update();
```
This form only executes the function if column `balance` has in fact changed value:

```sql
CREATE TRIGGER check_update
    BEFORE UPDATE ON accounts
    FOR EACH ROW
    WHEN (OLD.balance IS DISTINCT FROM NEW.balance)
    EXECUTE PROCEDURE check_account_update();
```

Call a function to log updates of `accounts`, but only if something changed:

```sql
CREATE TRIGGER log_update
    AFTER UPDATE ON accounts
    FOR EACH ROW
    WHEN (OLD.* IS DISTINCT FROM NEW.*)
    EXECUTE PROCEDURE log_account_update();
```

Execute the function `view_insert_row` for each row to insert rows into the tables underlying a view:

```sql
CREATE TRIGGER view_insert
    INSTEAD OF INSERT ON my_view
    FOR EACH ROW
    EXECUTE PROCEDURE view_insert_row();
```

Section 37.4 contains a complete example of a trigger function written in C.

**Compatibility**

The `CREATE TRIGGER` statement in PostgreSQL implements a subset of the SQL standard. The following functionalities are currently missing:

- SQL allows you to define aliases for the “old” and “new” rows or tables for use in the definition of the triggered action (e.g., `CREATE TRIGGER ... ON tablename REFERENCING OLD ROW AS somename NEW ROW AS othername ...`). Since PostgreSQL allows trigger procedures to be written in any number of user-defined languages, access to the data is handled in a language-specific way.
- PostgreSQL does not allow the old and new tables to be referenced in statement-level triggers, i.e., the tables that contain all the old and/or new rows, which are referred to by the `OLD TABLE` and `NEW TABLE` clauses in the SQL standard.
- PostgreSQL only allows the execution of a user-defined function for the triggered action. The standard allows the execution of a number of other SQL commands, such as `CREATE TABLE`, as the triggered action. This limitation is not hard to work around by creating a user-defined function that executes the desired commands.

SQL specifies that multiple triggers should be fired in time-of-creation order. PostgreSQL uses name order, which was judged to be more convenient.

SQL specifies that `BEFORE DELETE` triggers on cascaded deletes fire after the cascaded `DELETE` completes. The PostgreSQL behavior is for `BEFORE DELETE` to always fire before the delete action, even a cascading one. This is considered more consistent. There is also nonstandard behavior if `BEFORE`
triggers modify rows or prevent updates during an update that is caused by a referential action. This
 can lead to constraint violations or stored data that does not honor the referential constraint.

The ability to specify multiple actions for a single trigger using OR is a PostgreSQL extension of the
SQL standard.

The ability to fire triggers for TRUNCATE is a PostgreSQL extension of the SQL standard, as is the
ability to define statement-level triggers on views.

CREATE CONSTRAINT TRIGGER is a PostgreSQL extension of the SQL standard.

See Also

ALTER TRIGGER, DROP TRIGGER, CREATE FUNCTION, SET CONSTRAINTS
CREATE TYPE

Name

CREATE TYPE — define a new data type

Synopsis

CREATE TYPE name AS
( [ attribute_name data_type [ COLLATE collation ] [, ... ] ] )

CREATE TYPE name AS ENUM
( [ 'label' [, ... ] ] )

CREATE TYPE name AS RANGE (
  SUBTYPE = subtype
  [ , SUBTYPE_OPCCLASS = subtype_operator_class ]
  [ , COLLATION = collation ]
  [ , CANONICAL = canonical_function ]
  [ , SUBTYPE_DIFF = subtype_diff_function ]
)

CREATE TYPE name ( 
  INPUT = input_function,
  OUTPUT = output_function
  [ , RECEIVE = receive_function ]
  [ , SEND = send_function ]
  [ , TYPMOD_IN = type_modifier_input_function ]
  [ , TYPMOD_OUT = type_modifier_output_function ]
  [ , ANALYZE = analyze_function ]
  [ , INTERNALLENGTH = { internallength | VARIABLE } ]
  [ , PASSEDBYVALUE ]
  [ , ALIGNMENT = alignment ]
  [ , STORAGE = storage ]
  [ , LIKE = like_type ]
  [ , CATEGORY = category ]
  [ , PREFERRED = preferred ]
  [ , DEFAULT = default ]
  [ , ELEMENT = element ]
  [ , DELIMITER = delimiter ]
  [ , COLLATABLE = collatable ]
)

CREATE TYPE name

Description

CREATE TYPE registers a new data type for use in the current database. The user who defines a type becomes its owner.

If a schema name is given then the type is created in the specified schema. Otherwise it is created in the current schema. The type name must be distinct from the name of any existing type or domain
in the same schema. (Because tables have associated data types, the type name must also be distinct from the name of any existing table in the same schema.)

There are five forms of `CREATE TYPE`, as shown in the syntax synopsis above. They respectively create a composite type, an enum type, a range type, a base type, or a shell type. The first four of these are discussed in turn below. A shell type is simply a placeholder for a type to be defined later; it is created by issuing `CREATE TYPE` with no parameters except for the type name. Shell types are needed as forward references when creating range types and base types, as discussed in those sections.

**Composite Types**

The first form of `CREATE TYPE` creates a composite type. The composite type is specified by a list of attribute names and data types. An attribute’s collation can be specified too, if its data type is collatable. A composite type is essentially the same as the row type of a table, but using `CREATE TYPE` avoids the need to create an actual table when all that is wanted is to define a type. A stand-alone composite type is useful, for example, as the argument or return type of a function.

To be able to create a composite type, you must have `USAGE` privilege on all attribute types.

**Enumerated Types**

The second form of `CREATE TYPE` creates an enumerated (enum) type, as described in Section 8.7. Enum types take a list of quoted labels, each of which must be less than `NAMEDATALEN` bytes long (64 bytes in a standard PostgreSQL build). (It is possible to create an enumerated type with zero labels, but such a type cannot be used to hold values before at least one label is added using `ALTER TYPE`.)

**Range Types**

The third form of `CREATE TYPE` creates a new range type, as described in Section 8.17.

The range type’s `subtype` can be any type with an associated b-tree operator class (to determine the ordering of values for the range type). Normally the subtype’s default b-tree operator class is used to determine ordering; to use a non-default operator class, specify its name with `subtype_opclass`. If the subtype is collatable, and you want to use a non-default collation in the range’s ordering, specify the desired collation with the `collation` option.

The optional `canonical` function must take one argument of the range type being defined, and return a value of the same type. This is used to convert range values to a canonical form, when applicable. See Section 8.17.8 for more information. Creating a `canonical` function is a bit tricky, since it must be defined before the range type can be declared. To do this, you must first create a shell type, which is a placeholder type that has no properties except a name and an owner. This is done by issuing the command `CREATE TYPE name`, with no additional parameters. Then the function can be declared using the shell type as argument and result, and finally the range type can be declared using the same name. This automatically replaces the shell type entry with a valid range type.

The optional `subtype_diff` function must take two values of the `subtype` type as argument, and return a double precision value representing the difference between the two given values. While this is optional, providing it allows much greater efficiency of GiST indexes on columns of the range type. See Section 8.17.8 for more information.
Base Types

The fourth form of `CREATE TYPE` creates a new base type (scalar type). To create a new base type, you must be a superuser. (This restriction is made because an erroneous type definition could confuse or even crash the server.)

The parameters can appear in any order, not only that illustrated above, and most are optional. You must register two or more functions (using `CREATE FUNCTION`) before defining the type. The support functions `input_function` and `output_function` are required, while the functions `receive_function`, `send_function`, `type_modifier_input_function`, `type_modifier_output_function` and `analyze_function` are optional. Generally these functions have to be coded in C or another low-level language.

The `input_function` converts the type’s external textual representation to the internal representation used by the operators and functions defined for the type. `output_function` performs the reverse transformation. The input function can be declared as taking one argument of type `cstring`, or as taking three arguments of types `cstring`, `oid`, `integer`. The first argument is the input text as a C string, the second argument is the type’s own OID (except for array types, which instead receive their element type’s OID), and the third is the `typmod` of the destination column, if known (-1 will be passed if not). The input function must return a value of the data type itself. Usually, an input function should be declared `STRICT`; if it is not, it will be called with a NULL first parameter when reading a NULL input value. The function must still return NULL in this case, unless it raises an error. (This case is mainly meant to support domain input functions, which might need to reject NULL inputs.) The output function must be declared as taking one argument of the new data type. The output function must return type `cstring`. Output functions are not invoked for NULL values.

The optional `receive_function` converts the type’s external binary representation to the internal representation. If this function is not supplied, the type cannot participate in binary input. The binary representation should be chosen to be cheap to convert to internal form, while being reasonably portable. (For example, the standard integer data types use network byte order as the external binary representation, while the internal representation is in the machine’s native byte order.) The receive function should perform adequate checking to ensure that the value is valid. The receive function can be declared as taking one argument of type `internal`, or as taking three arguments of types `internal`, `oid`, `integer`. The first argument is a pointer to a `StringInfo` buffer holding the received byte string; the optional arguments are the same as for the text input function. The receive function must return a value of the data type itself. Usually, a receive function should be declared `STRICT`; if it is not, it will be called with a NULL first parameter when reading a NULL input value. The function must still return NULL in this case, unless it raises an error. (This case is mainly meant to support domain receive functions, which might need to reject NULL inputs.) Similarly, the optional `send_function` converts from the internal representation to the external binary representation. If this function is not supplied, the type cannot participate in binary output. The send function must be declared as taking one argument of the new data type. The send function must return type `bytea`. Send functions are not invoked for NULL values.

You should at this point be wondering how the input and output functions can be declared to have results or arguments of the new type, when they have to be created before the new type can be created. The answer is that the type should first be defined as a *shell type*, which is a placeholder type that has no properties except a name and an owner. This is done by issuing the command `CREATE TYPE name`, with no additional parameters. Then the C I/O functions can be defined referencing the shell type. Finally, `CREATE TYPE` with a full definition replaces the shell entry with a complete, valid type definition, after which the new type can be used normally.

The optional `type_modifier_input_function` and `type_modifier_output_function` are needed if the type supports modifiers, that is optional constraints attached to a type declaration, such as `char(5)` or `numeric(30,2)`. PostgreSQL allows user-defined types to take one or
CREATE TYPE

more simple constants or identifiers as modifiers. However, this information must be capable of being packed into a single non-negative integer value for storage in the system catalogs. The type_modifier_input_function is passed the declared modifier(s) in the form of a cstring array. It must check the values for validity (throwing an error if they are wrong), and if they are correct, return a single non-negative integer value that will be stored as the column “typmod”. Type modifiers will be rejected if the type does not have a type_modifier_input_function. The type_modifier_output_function converts the internal integer typmod value back to the correct form for user display. It must return a cstring value that is the exact string to append to the type name; for example numeric’s function might return (30,2). It is allowed to omit the type_modifier_output_function, in which case the default display format is just the stored typmod integer value enclosed in parentheses.

The optional analyze_function performs type-specific statistics collection for columns of the data type. By default, ANALYZE will attempt to gather statistics using the type’s “equals” and “less-than” operators, if there is a default b-tree operator class for the type. For non-scalar types this behavior is likely to be unsuitable, so it can be overridden by specifying a custom analysis function. The analysis function must be declared to take a single argument of type internal, and return a boolean result. The detailed API for analysis functions appears in src/include/commands/vacuum.h.

While the details of the new type’s internal representation are only known to the I/O functions and other functions you create to work with the type, there are several properties of the internal representation that must be declared to PostgreSQL. Foremost of these is internallength. Base data types can be fixed-length, in which case internallength is a positive integer, or variable-length, indicated by setting internallength to VARIABLE. (Internally, this is represented by setting typlen to -1.) The internal representation of all variable-length types must start with a 4-byte integer giving the total length of this value of the type. (Note that the length field is often encoded, as described in Section 65.2; it’s unwise to access it directly.)

The optional flag PASSEDBYVALUE indicates that values of this data type are passed by value, rather than by reference. Types passed by value must be fixed-length, and their internal representation cannot be larger than the size of the Datum type (4 bytes on some machines, 8 bytes on others).

The alignment parameter specifies the storage alignment required for the data type. The allowed values equate to alignment on 1, 2, 4, or 8 byte boundaries. Note that variable-length types must have an alignment of at least 4, since they necessarily contain an int4 as their first component.

The storage parameter allows selection of storage strategies for variable-length data types. (Only plain is allowed for fixed-length types.) plain specifies that data of the type will always be stored in-line and not compressed. extended specifies that the system will first try to compress a long data value, and will move the value out of the main table row if it’s still too long. external allows the value to be moved out of the main table, but the system will not try to compress it. main allows compression, but discourages moving the value out of the main table. (Data items with this storage strategy might still be moved out of the main table if there is no other way to make a row fit, but they will be kept in the main table preferentially over extended and external items.)

All storage values other than plain imply that the functions of the data type can handle values that have been toasted, as described in Section 65.2 and Section 36.11.1. The specific other value given merely determines the default TOAST storage strategy for columns of a toastable data type; users can pick other strategies for individual columns using ALTER TABLE SET STORAGE.

The like_type parameter provides an alternative method for specifying the basic representation properties of a data type: copy them from some existing type. The values of internallength, passedbyvalue, alignment, and storage are copied from the named type. (It is possible, though usually undesirable, to override some of these values by specifying them along with the LIKE clause.) Specifying representation this way is especially useful when the low-level implementation of the new type “piggybacks” on an existing type in some fashion.
CREATE TYPE

The category and preferred parameters can be used to help control which implicit cast will be applied in ambiguous situations. Each data type belongs to a category named by a single ASCII character, and each type is either “preferred” or not within its category. The parser will prefer casting to preferred types (but only from other types within the same category) when this rule is helpful in resolving overloaded functions or operators. For more details see Chapter 10. For types that have no implicit casts to or from any other types, it is sufficient to leave these settings at the defaults. However, for a group of related types that have implicit casts, it is often helpful to mark them all as belonging to a category and select one or two of the “most general” types as being preferred within the category. The category parameter is especially useful when adding a user-defined type to an existing built-in category, such as the numeric or string types. However, it is also possible to create new entirely-user-defined type categories. Select any ASCII character other than an upper-case letter to name such a category.

A default value can be specified, in case a user wants columns of the data type to default to something other than the null value. Specify the default with the DEFAULT key word. (Such a default can be overridden by an explicit DEFAULT clause attached to a particular column.)

To indicate that a type is an array, specify the type of the array elements using the ELEMENT key word. For example, to define an array of 4-byte integers (int4), specify ELEMENT = int4. More details about array types appear below.

To indicate the delimiter to be used between values in the external representation of arrays of this type, delimiter can be set to a specific character. The default delimiter is the comma (,). Note that the delimiter is associated with the array element type, not the array type itself.

If the optional Boolean parameter collatable is true, column definitions and expressions of the type may carry collation information through use of the COLLATE clause. It is up to the implementations of the functions operating on the type to actually make use of the collation information; this does not happen automatically merely by marking the type collatable.

Array Types

Whenever a user-defined type is created, PostgreSQL automatically creates an associated array type, whose name consists of the element type’s name prepended with an underscore, and truncated if necessary to keep it less than NAMEDATALEN bytes long. (If the name so generated collides with an existing type name, the process is repeated until a non-colliding name is found.) This implicitly-created array type is variable length and uses the built-in input and output functions array_in and array_out. The array type tracks any changes in its element type’s owner or schema, and is dropped if the element type is.

You might reasonably ask why there is an ELEMENT option, if the system makes the correct array type automatically. The only case where it’s useful to use ELEMENT is when you are making a fixed-length type that happens to be internally an array of a number of identical things, and you want to allow these things to be accessed directly by subscripting, in addition to whatever operations you plan to provide for the type as a whole. For example, type point is represented as just two floating-point numbers, which can be accessed using point[0] and point[1]. Note that this facility only works for fixed-length types whose internal form is exactly a sequence of identical fixed-length fields. A subscriptable variable-length type must have the generalized internal representation used by array_in and array_out. For historical reasons (i.e., this is clearly wrong but it’s far too late to change it), subscripting of fixed-length array types starts from zero, rather than from one as for variable-length arrays.
CREATE TYPE

Parameters

**name**

The name (optionally schema-qualified) of a type to be created.

**attribute_name**

The name of an attribute (column) for the composite type.

**data_type**

The name of an existing data type to become a column of the composite type.

**collation**

The name of an existing collation to be associated with a column of a composite type, or with a range type.

**label**

A string literal representing the textual label associated with one value of an enum type.

**subtype**

The name of the element type that the range type will represent ranges of.

**subtype_operator_class**

The name of a b-tree operator class for the subtype.

**canonical_function**

The name of the canonicalization function for the range type.

**subtype_diff_function**

The name of a difference function for the subtype.

**input_function**

The name of a function that converts data from the type’s external textual form to its internal form.

**output_function**

The name of a function that converts data from the type’s internal form to its external textual form.

**receive_function**

The name of a function that converts data from the type’s external binary form to its internal form.

**send_function**

The name of a function that converts data from the type’s internal form to its external binary form.

**type_modifier_input_function**

The name of a function that converts an array of modifier(s) for the type into internal form.

**type_modifier_output_function**

The name of a function that converts the internal form of the type’s modifier(s) to external textual form.

**analyze_function**

The name of a function that performs statistical analysis for the data type.
CREATE TYPE

`internallength`  
A numeric constant that specifies the length in bytes of the new type’s internal representation. The default assumption is that it is variable-length.

`alignment`  
The storage alignment requirement of the data type. If specified, it must be `char`, `int2`, `int4`, or `double`; the default is `int4`.

`storage`  
The storage strategy for the data type. If specified, must be `plain`, `external`, `extended`, or `main`; the default is `plain`.

`like_type`  
The name of an existing data type that the new type will have the same representation as. The values of `internallength`, `passedbyvalue`, `alignment`, and `storage` are copied from that type, unless overridden by explicit specification elsewhere in this `CREATE TYPE` command.

`category`  
The category code (a single ASCII character) for this type. The default is `'U'` for “user-defined type”. Other standard category codes can be found in Table 50-56. You may also choose other ASCII characters in order to create custom categories.

`preferred`  
True if this type is a preferred type within its type category, else false. The default is false. Be very careful about creating a new preferred type within an existing type category, as this could cause surprising changes in behavior.

`default`  
The default value for the data type. If this is omitted, the default is null.

`element`  
The type being created is an array; this specifies the type of the array elements.

`delimiter`  
The delimiter character to be used between values in arrays made of this type.

`collatable`  
True if this type’s operations can use collation information. The default is false.

**Notes**

Because there are no restrictions on use of a data type once it’s been created, creating a base type or range type is tantamount to granting public execute permission on the functions mentioned in the type definition. This is usually not an issue for the sorts of functions that are useful in a type definition. But you might want to think twice before designing a type in a way that would require “secret” information to be used while converting it to or from external form.

Before PostgreSQL version 8.3, the name of a generated array type was always exactly the element type’s name with one underscore character (`_`) prepended. (Type names were therefore restricted in length to one less character than other names.) While this is still usually the case, the array type name may vary from this in case of maximum-length names or collisions with user type names that begin with underscore. Writing code that depends on this convention is therefore deprecated. Instead, use `pg_type.typarray` to locate the array type associated with a given type.
CREATE TYPE

It may be advisable to avoid using type and table names that begin with underscore. While the server will change generated array type names to avoid collisions with user-given names, there is still risk of confusion, particularly with old client software that may assume that type names beginning with underscores always represent arrays.

Before PostgreSQL version 8.2, the shell-type creation syntax `CREATE TYPE` name did not exist. The way to create a new base type was to create its input function first. In this approach, PostgreSQL will first see the name of the new data type as the return type of the input function. The shell type is implicitly created in this situation, and then it can be referenced in the definitions of the remaining I/O functions. This approach still works, but is deprecated and might be disallowed in some future release. Also, to avoid accidentally cluttering the catalogs with shell types as a result of simple typos in function definitions, a shell type will only be made this way when the input function is written in C.

In PostgreSQL versions before 7.3, it was customary to avoid creating a shell type at all, by replacing the functions’ forward references to the type name with the placeholder pseudotype `opaque`. The `cstring` arguments and results also had to be declared as `opaque` before 7.3. To support loading of old dump files, `CREATE TYPE` will accept I/O functions declared using `opaque`, but it will issue a notice and change the function declarations to use the correct types.

**Examples**

This example creates a composite type and uses it in a function definition:

```sql
CREATE TYPE compfoo AS (f1 int, f2 text);
CREATE FUNCTION getfoo() RETURNS SETOF compfoo AS $$
    SELECT fooid, fooname FROM foo
$$ LANGUAGE SQL;
```

This example creates an enumerated type and uses it in a table definition:

```sql
CREATE TYPE bug_status AS ENUM ('new', 'open', 'closed');
CREATE TABLE bug (id serial, description text, status bug_status);
```

This example creates a range type:

```sql
CREATE TYPE float8_range AS RANGE (subtype = float8, subtype_diff = float8mi);
```

This example creates the base data type `box` and then uses the type in a table definition:

```sql
CREATE TYPE box;
CREATE FUNCTION my_box_in_function(cstring) RETURNS box AS ... ;
CREATE FUNCTION my_box_out_function(box) RETURNS cstring AS ... ;
```
CREATE TYPE box (  
  INTERNALLENGTH = 16,  
  INPUT = my_box_in_function,  
  OUTPUT = my_box_out_function  
);  

CREATE TABLE myboxes (  
  id integer,  
  description box  
);  

If the internal structure of box were an array of four float4 elements, we might instead use:  

CREATE TYPE box (  
  INTERNALLENGTH = 16,  
  INPUT = my_box_in_function,  
  OUTPUT = my_box_out_function,  
  ELEMENT = float4  
);  

which would allow a box value’s component numbers to be accessed by subscripting. Otherwise the type behaves the same as before.  

This example creates a large object type and uses it in a table definition:  

CREATE TYPE bigobj (  
  INPUT = lo_filein, OUTPUT = lo_fileout,  
  INTERNALLENGTH = VARIABLE  
);  
CREATE TABLE big_objs (  
  id integer,  
  obj bigobj  
);  

More examples, including suitable input and output functions, are in Section 36.11.  

Compatibility  

The first form of the CREATE TYPE command, which creates a composite type, conforms to the SQL standard. The other forms are PostgreSQL extensions. The CREATE TYPE statement in the SQL standard also defines other forms that are not implemented in PostgreSQL.  

The ability to create a composite type with zero attributes is a PostgreSQL-specific deviation from the standard (analogous to the same case in CREATE TABLE).  

See Also  

ALTER TYPE, CREATE DOMAIN, CREATE FUNCTION, DROP TYPE
CREATE USER

Name

CREATE USER — define a new database role

Synopsis

CREATE USER name [ [ WITH ] option [ ... ] ]

where option can be:

    SUPERUSER | NOSUPERUSER
| CREATEDB | NOCREATEDB
| CREATEROLE | NOCREATEROLE
| INHERIT | NOINHERIT
| LOGIN | NOLOGIN
| REPLICATION | NOREPLICATION
| BYPASSRLS | NOBYPASSRLS
| CONNECTION LIMIT connlimit
| [ ENCRYPTED | UNENCRYPTED ] PASSWORD 'password'
| VALID UNTIL 'timestamp'
| IN ROLE role_name [, ...]
| IN GROUP role_name [, ...]
| ROLE role_name [, ...]
| ADMIN role_name [, ...]
| USER role_name [, ...]
| SYSID uid

Description

CREATE USER is now an alias for CREATE ROLE. The only difference is that when the command is spelled CREATE USER, LOGIN is assumed by default, whereas NOLOGIN is assumed when the command is spelled CREATE ROLE.

Compatibility

The CREATE USER statement is a PostgreSQL extension. The SQL standard leaves the definition of users to the implementation.

See Also

CREATE ROLE
CREATE USER MAPPING

Name
CREATE USER MAPPING — define a new mapping of a user to a foreign server

Synopsis
CREATE USER MAPPING FOR { user_name | USER | CURRENT_USER | PUBLIC }
   SERVER server_name
   [ OPTIONS ( option 'value' [ , ... ] ) ]

Description
CREATE USER MAPPING defines a mapping of a user to a foreign server. A user mapping typically encapsulates connection information that a foreign-data wrapper uses together with the information encapsulated by a foreign server to access an external data resource.

The owner of a foreign server can create user mappings for that server for any user. Also, a user can create a user mapping for their own user name if USAGE privilege on the server has been granted to the user.

Parameters
user_name

The name of an existing user that is mapped to foreign server. CURRENT_USER and USER match the name of the current user. When PUBLIC is specified, a so-called public mapping is created that is used when no user-specific mapping is applicable.

server_name

The name of an existing server for which the user mapping is to be created.

OPTIONS ( option 'value' [ , ... ] )

This clause specifies the options of the user mapping. The options typically define the actual user name and password of the mapping. Option names must be unique. The allowed option names and values are specific to the server’s foreign-data wrapper.

Examples
Create a user mapping for user bob, server foo:

CREATE USER MAPPING FOR bob SERVER foo OPTIONS (user 'bob', password 'secret');
Compatibility

CREATE USER MAPPING conforms to ISO/IEC 9075-9 (SQL/MED).

See Also

ALTER USER MAPPING, DROP USER MAPPING, CREATE FOREIGN DATA WRAPPER, CREATE SERVER
CREATE VIEW

**Name**

CREATE VIEW — define a new view

**Synopsis**

```
CREATE [ OR REPLACE ] [ TEMP | TEMPORARY ] [ RECURSIVE ] VIEW name [ ( column_name [, ] ... ) ]
[ WITH ( view_option_name = view_option_value [, ] ... ) ]
AS query
[ WITH [ CASCADED | LOCAL ] CHECK OPTION ]
```

**Description**

CREATE VIEW defines a view of a query. The view is not physically materialized. Instead, the query is run every time the view is referenced in a query.

CREATE OR REPLACE VIEW is similar, but if a view of the same name already exists, it is replaced. The new query must generate the same columns that were generated by the existing view query (that is, the same column names in the same order and with the same data types), but it may add additional columns to the end of the list. The calculations giving rise to the output columns may be completely different.

If a schema name is given (for example, CREATE VIEW myschema.myview ...) then the view is created in the specified schema. Otherwise it is created in the current schema. Temporary views exist in a special schema, so a schema name cannot be given when creating a temporary view. The name of the view must be distinct from the name of any other view, table, sequence, index or foreign table in the same schema.

**Parameters**

**TEMPORARY or TEMP**

If specified, the view is created as a temporary view. Temporary views are automatically dropped at the end of the current session. Existing permanent relations with the same name are not visible to the current session while the temporary view exists, unless they are referenced with schema-qualified names.

If any of the tables referenced by the view are temporary, the view is created as a temporary view (whether TEMPORARY is specified or not).

**RECURSIVE**

Creates a recursive view. The syntax

```
CREATE RECURSIVE VIEW [ schema . ] view_name (column_names) AS SELECT ...;
```

is equivalent to

```
CREATE VIEW [ schema . ] view_name AS WITH RECURSIVE view_name (column_names) AS (SELECT ...);
```

A view column name list must be specified for a recursive view.
CREATE VIEW

name

The name (optionally schema-qualified) of a view to be created.

column_name

An optional list of names to be used for columns of the view. If not given, the column names are deduced from the query.

WITH ( view_option_name [= view_option_value] [, ... ] )

This clause specifies optional parameters for a view; the following parameters are supported:

check_option (string)

This parameter may be either local or cascaded, and is equivalent to specifying WITH [ CASCADED | LOCAL ] CHECK OPTION (see below). This option can be changed on existing views using ALTER VIEW.

security_barrier (boolean)

This should be used if the view is intended to provide row-level security. See Section 39.5 for full details.

query

A SELECT or VALUES command which will provide the columns and rows of the view.

WITH [ CASCADED | LOCAL ] CHECK OPTION

This option controls the behavior of automatically updatable views. When this option is specified, INSERT and UPDATE commands on the view will be checked to ensure that new rows satisfy the view-defining condition (that is, the new rows are checked to ensure that they are visible through the view). If they are not, the update will be rejected. If the CHECK OPTION is not specified, INSERT and UPDATE commands on the view are allowed to create rows that are not visible through the view. The following check options are supported:

LOCAL

New rows are only checked against the conditions defined directly in the view itself. Any conditions defined on underlying base views are not checked (unless they also specify the CHECK OPTION).

CASCADED

New rows are checked against the conditions of the view and all underlying base views. If the CHECK OPTION is specified, and neither LOCAL nor CASCADED is specified, then CASCADED is assumed.

The CHECK OPTION may not be used with RECURSIVE views.

Note that the CHECK OPTION is only supported on views that are automatically updatable, and do not have INSTEAD OF triggers or INSTEAD rules. If an automatically updatable view is defined on top of a base view that has INSTEAD OF triggers, then the LOCAL CHECK OPTION may be used to check the conditions on the automatically updatable view, but the conditions on the base view with INSTEAD OF triggers will not be checked (a cascaded check option will not cascade down to a trigger-updatable view, and any check options defined directly on a trigger-updatable view will be ignored). If the view or any of its base relations has an INSTEAD rule that causes the INSERT or UPDATE command to be rewritten, then all check options will be ignored in the
CREATE VIEW

rewritten query, including any checks from automatically updatable views defined on top of the relation with the INSTEAD rule.

Notes

Use the DROP VIEW statement to drop views.

Be careful that the names and types of the view’s columns will be assigned the way you want. For example:

CREATE VIEW vista AS SELECT ‘Hello World’;

is bad form in two ways: the column name defaults to column, and the column data type defaults to unknown. If you want a string literal in a view’s result, use something like:

CREATE VIEW vista AS SELECT text ‘Hello World’ AS hello;

Access to tables referenced in the view is determined by permissions of the view owner. In some cases, this can be used to provide secure but restricted access to the underlying tables. However, not all views are secure against tampering; see Section 39.5 for details. Functions called in the view are treated the same as if they had been called directly from the query using the view. Therefore the user of a view must have permissions to call all functions used by the view.

When CREATE OR REPLACE VIEW is used on an existing view, only the view’s defining SELECT rule is changed. Other view properties, including ownership, permissions, and non-SELECT rules, remain unchanged. You must own the view to replace it (this includes being a member of the owning role).

Updatable Views

Simple views are automatically updatable: the system will allow INSERT, UPDATE and DELETE statements to be used on the view in the same way as on a regular table. A view is automatically updatable if it satisfies all of the following conditions:

- The view must have exactly one entry in its FROM list, which must be a table or another updatable view.
- The view definition must not contain WITH, DISTINCT, GROUP BY, HAVING, LIMIT, or OFFSET clauses at the top level.
- The view definition must not contain set operations (UNION, INTERSECT or EXCEPT) at the top level.
- The view’s select list must not contain any aggregates, window functions or set-returning functions.

An automatically updatable view may contain a mix of updatable and non-updatable columns. A column is updatable if it is a simple reference to an updatable column of the underlying base relation; otherwise the column is read-only, and an error will be raised if an INSERT or UPDATE statement attempts to assign a value to it.

If the view is automatically updatable the system will convert any INSERT, UPDATE or DELETE statement on the view into the corresponding statement on the underlying base relation. INSERT statements that have an ON CONFLICT UPDATE clause are fully supported.
CREATE VIEW

If an automatically updatable view contains a WHERE condition, the condition restricts which rows of the base relation are available to be modified by UPDATE and DELETE statements on the view. However, an UPDATE is allowed to change a row so that it no longer satisfies the WHERE condition, and thus is no longer visible through the view. Similarly, an INSERT command can potentially insert base-relation rows that do not satisfy the WHERE condition and thus are not visible through the view. The CHECK OPTION may be used to prevent INSERT and UPDATE commands from creating such rows that are not visible through the view.

If an automatically updatable view is marked with the security_barrier property then all the view’s WHERE conditions (and any conditions using operators which are marked as LEAKPROOF) will always be evaluated before any conditions that a user of the view has added. See Section 39.5 for full details. Note that, due to this, rows which are not ultimately returned (because they do not pass the user’s WHERE conditions) may still end up being locked. EXPLAIN can be used to see which conditions are applied at the relation level (and therefore do not lock rows) and which are not.

A more complex view that does not satisfy all these conditions is read-only by default: the system will not allow an insert, update, or delete on the view. You can get the effect of an updatable view by creating INSTEAD OF triggers on the view, which must convert attempted inserts, etc. on the view into appropriate actions on other tables. For more information see CREATE TRIGGER. Another possibility is to create rules (see CREATE RULE), but in practice triggers are easier to understand and use correctly.

Note that the user performing the insert, update or delete on the view must have the corresponding insert, update or delete privilege on the view. In addition the view’s owner must have the relevant privileges on the underlying base relations, but the user performing the update does not need any permissions on the underlying base relations (see Section 39.5).

Examples

Create a view consisting of all comedy films:

```
CREATE VIEW comedies AS
    SELECT *
    FROM films
    WHERE kind = 'Comedy';
```

This will create a view containing the columns that are in the film table at the time of view creation. Though * was used to create the view, columns added later to the table will not be part of the view.

Create a view with LOCAL CHECK OPTION:

```
CREATE VIEW universal_comedies AS
    SELECT *
    FROM comedies
    WHERE classification = 'U'
    WITH LOCAL CHECK OPTION;
```

This will create a view based on the comedies view, showing only films with kind = 'Comedy' and classification = 'U'. Any attempt to INSERT or UPDATE a row in the view will be rejected if the new row doesn’t have classification = 'U', but the film kind will not be checked.

Create a view with CASCADED CHECK OPTION:

```
CREATE VIEW pg_comedies AS
```
CREATE VIEW

SELECT *
FROM comedies
WHERE classification = 'PG'
WITH CASCADED CHECK OPTION;

This will create a view that checks both the kind and classification of new rows.

Create a view with a mix of updatable and non-updatable columns:

CREATE VIEW comedies AS
SELECT f.*,
  country_code_to_name(f.country_code) AS country,
  (SELECT avg(r.rating)
    FROM user_ratings r
    WHERE r.film_id = f.id) AS avg_rating
FROM films f
WHERE f.kind = 'Comedy';

This view will support INSERT, UPDATE and DELETE. All the columns from the films table will be updatable, whereas the computed columns country and avg_rating will be read-only.

Create a recursive view consisting of the numbers from 1 to 100:

CREATE RECURSIVE VIEW public.nums_1_100 (n) AS
VALUES (1)
UNION ALL
  SELECT n+1 FROM nums_1_100 WHERE n < 100;

Notice that although the recursive view’s name is schema-qualified in this CREATE, its internal self-reference is not schema-qualified. This is because the implicitly-created CTE’s name cannot be schema-qualified.

Compatibility

CREATE OR REPLACE VIEW is a PostgreSQL language extension. So is the concept of a temporary view. The WITH ( ... ) clause is an extension as well.

See Also

ALTER VIEW, DROP VIEW, CREATE MATERIALIZED VIEW
DEALLOCATE

Name
DEALLOCATE — deallocate a prepared statement

Synopsis
DEALLOCATE [ PREPARE ] { name | ALL }

Description
DEALLOCATE is used to deallocate a previously prepared SQL statement. If you do not explicitly deallocate a prepared statement, it is deallocated when the session ends.
For more information on prepared statements, see PREPARE.

Parameters

PREPARE
This key word is ignored.

name
The name of the prepared statement to deallocate.

ALL
Deallocate all prepared statements.

Compatibility
The SQL standard includes a DEALLOCATE statement, but it is only for use in embedded SQL.

See Also
EXECUTE, PREPARE
DECLARE

Name

DECLARE — define a cursor

Synopsis

DECLARE name [ BINARY ] [ INSENSITIVE ] [ [ NO ] SCROLL ]
CURSOR [ [ WITH | WITHOUT ] HOLD ] FOR query

Description

DECLARE allows a user to create cursors, which can be used to retrieve a small number of rows at a time out of a larger query. After the cursor is created, rows are fetched from it using FETCH.

Note: This page describes usage of cursors at the SQL command level. If you are trying to use cursors inside a PL/pgSQL function, the rules are different — see Section 41.7.

Parameters

name

The name of the cursor to be created.

BINARY

Causes the cursor to return data in binary rather than in text format.

INSENSITIVE

Indicates that data retrieved from the cursor should be unaffected by updates to the table(s) underlying the cursor that occur after the cursor is created. In PostgreSQL, this is the default behavior; so this key word has no effect and is only accepted for compatibility with the SQL standard.

SCROLL

NO SCROLL

SCROLL specifies that the cursor can be used to retrieve rows in a nonsequential fashion (e.g., backward). Depending upon the complexity of the query's execution plan, specifying SCROLL might impose a performance penalty on the query's execution time. NO SCROLL specifies that the cursor cannot be used to retrieve rows in a nonsequential fashion. The default is to allow scrolling in some cases; this is not the same as specifying SCROLL. See Notes for details.

WITH HOLD

WITHOUT HOLD

WITH HOLD specifies that the cursor can continue to be used after the transaction that created it successfully commits. WITHOUT HOLD specifies that the cursor cannot be used outside of the
transaction that created it. If neither WITHOUT HOLD nor WITH HOLD is specified, WITHOUT HOLD is the default.

query

A SELECT or VALUES command which will provide the rows to be returned by the cursor.

The key words BINARY, INSENSITIVE, and SCROLL can appear in any order.

Notes

Normal cursors return data in text format, the same as a SELECT would produce. The BINARY option specifies that the cursor should return data in binary format. This reduces conversion effort for both the server and client, at the cost of more programmer effort to deal with platform-dependent binary data formats. As an example, if a query returns a value of one from an integer column, you would get a string of 1 with a default cursor, whereas with a binary cursor you would get a 4-byte field containing the internal representation of the value (in big-endian byte order).

Binary cursors should be used carefully. Many applications, including psql, are not prepared to handle binary cursors and expect data to come back in the text format.

Note: When the client application uses the "extended query" protocol to issue a `FETCH` command, the Bind protocol message specifies whether data is to be retrieved in text or binary format. This choice overrides the way that the cursor is defined. The concept of a binary cursor as such is thus obsolete when using extended query protocol — any cursor can be treated as either text or binary.

Unless WITH HOLD is specified, the cursor created by this command can only be used within the current transaction. Thus, DECLARE without WITH HOLD is useless outside a transaction block: the cursor would survive only to the completion of the statement. Therefore PostgreSQL reports an error if such a command is used outside a transaction block. Use BEGIN and COMMIT (or ROLLBACK) to define a transaction block.

If WITH HOLD is specified and the transaction that created the cursor successfully commits, the cursor can continue to be accessed by subsequent transactions in the same session. (But if the creating transaction is aborted, the cursor is removed.) A cursor created with WITH HOLD is closed when an explicit CLOSE command is issued on it, or the session ends. In the current implementation, the rows represented by a held cursor are copied into a temporary file or memory area so that they remain available for subsequent transactions.

WITH HOLD may not be specified when the query includes FOR UPDATE or FOR SHARE.

The SCROLL option should be specified when defining a cursor that will be used to fetch backwards. This is required by the SQL standard. However, for compatibility with earlier versions, PostgreSQL will allow backward fetches without SCROLL, if the cursor’s query plan is simple enough that no extra overhead is needed to support it. However, application developers are advised not to rely on using backward fetches from a cursor that has not been created with SCROLL. If NO SCROLL is specified, then backward fetches are disallowed in any case.

Backward fetches are also disallowed when the query includes FOR UPDATE or FOR SHARE; therefore SCROLL may not be specified in this case.
DECLARE

Caution
Scorable and WITH HOLD cursors may give unexpected results if they invoke any volatile functions (see Section 36.6). When a previously fetched row is re-fetched, the functions might be re-executed, perhaps leading to results different from the first time. One workaround for such cases is to declare the cursor WITH HOLD and commit the transaction before reading any rows from it. This will force the entire output of the cursor to be materialized in temporary storage, so that volatile functions are executed exactly once for each row.

If the cursor’s query includes FOR UPDATE or FOR SHARE, then returned rows are locked at the time they are first fetched, in the same way as for a regular SELECT command with these options. In addition, the returned rows will be the most up-to-date versions; therefore these options provide the equivalent of what the SQL standard calls a “sensitive cursor”. (Specifying INSENSITIVE together with FOR UPDATE or FOR SHARE is an error.)

Caution
It is generally recommended to use FOR UPDATE if the cursor is intended to be used with UPDATE ... WHERE CURRENT OF or DELETE ... WHERE CURRENT OF. Using FOR UPDATE prevents other sessions from changing the rows between the time they are fetched and the time they are updated. Without FOR UPDATE, a subsequent WHERE CURRENT OF command will have no effect if the row was changed since the cursor was created.

Another reason to use FOR UPDATE is that without it, a subsequent WHERE CURRENT OF might fail if the cursor query does not meet the SQL standard's rules for being “simply updatable” (in particular, the cursor must reference just one table and not use grouping or ORDER BY). Cursors that are not simply updatable might work, or might not, depending on plan choice details; so in the worst case, an application might work in testing and then fail in production. If FOR UPDATE is specified, the cursor is guaranteed to be updatable.

The main reason not to use FOR UPDATE with WHERE CURRENT OF is if you need the cursor to be scrollable, or to be insensitive to the subsequent updates (that is, continue to show the old data). If this is a requirement, pay close heed to the caveats shown above.

The SQL standard only makes provisions for cursors in embedded SQL. The PostgreSQL server does not implement an OPEN statement for cursors; a cursor is considered to be open when it is declared. However, ECPG, the embedded SQL preprocessor for PostgreSQL, supports the standard SQL cursor conventions, including those involving DECLARE and OPEN statements.

You can see all available cursors by querying the pg_cursors system view.

Examples
To declare a cursor:

DECLARE liahona CURSOR FOR SELECT * FROM films;

See FETCH for more examples of cursor usage.
Compatibility

The SQL standard says that it is implementation-dependent whether cursors are sensitive to concurrent updates of the underlying data by default. In PostgreSQL, cursors are insensitive by default, and can be made sensitive by specifying `FOR UPDATE`. Other products may work differently.

The SQL standard allows cursors only in embedded SQL and in modules. PostgreSQL permits cursors to be used interactively.

Binary cursors are a PostgreSQL extension.

See Also

CLOSE, FETCH, MOVE
DELETE

Name

DELETE — delete rows of a table

Synopsis

[ WITH [ RECURSIVE ] with_query [, ...] ]
DELETE FROM [ ONLY ] table_name [ * ] [ [ AS ] alias ]
[ USING using_list ]
[ WHERE condition ] WHERE CURRENT OF cursor_name ]
[ RETURNING * | output_expression [ [ AS ] output_name ] [, ...] ]

Description

DELETE deletes rows that satisfy the WHERE clause from the specified table. If the WHERE clause is absent, the effect is to delete all rows in the table. The result is a valid, but empty table.

Tip: TRUNCATE is a PostgreSQL extension that provides a faster mechanism to remove all rows from a table.

There are two ways to delete rows in a table using information contained in other tables in the database: using sub-selects, or specifying additional tables in the USING clause. Which technique is more appropriate depends on the specific circumstances.

The optional RETURNING clause causes DELETE to compute and return value(s) based on each row actually deleted. Any expression using the table’s columns, and/or columns of other tables mentioned in USING, can be computed. The syntax of the RETURNING list is identical to that of the output list of SELECT.

You must have the DELETE privilege on the table to delete from it, as well as the SELECT privilege for any table in the USING clause or whose values are read in the condition.

Parameters

with_query

The WITH clause allows you to specify one or more subqueries that can be referenced by name in the DELETE query. See Section 7.8 and SELECT for details.

table_name

The name (optionally schema-qualified) of the table to delete rows from. If ONLY is specified before the table name, matching rows are deleted from the named table only. If ONLY is not specified, matching rows are also deleted from any tables inheriting from the named table. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.
DELETE

alias

A substitute name for the target table. When an alias is provided, it completely hides the actual name of the table. For example, given `DELETE FROM foo AS f`, the remainder of the DELETE statement must refer to this table as `f` not `foo`.

using_list

A list of table expressions, allowing columns from other tables to appear in the WHERE condition. This is similar to the list of tables that can be specified in the FROM Clause of a SELECT statement; for example, an alias for the table name can be specified. Do not repeat the target table in the using_list, unless you wish to set up a self-join.

condition

An expression that returns a value of type boolean. Only rows for which this expression returns true will be deleted.

cursor_name

The name of the cursor to use in a WHERE CURRENT OF condition. The row to be deleted is the one most recently fetched from this cursor. The cursor must be a non-grouping query on the DELETE’s target table. Note that WHERE CURRENT OF cannot be specified together with a Boolean condition. See DECLARE for more information about using cursors with WHERE CURRENT OF.

output_expression

An expression to be computed and returned by the DELETE command after each row is deleted. The expression can use any column names of the table named by table_name or table(s) listed in USING. Write * to return all columns.

output_name

A name to use for a returned column.

Outputs

On successful completion, a DELETE command returns a command tag of the form

```
DELETE count
```

The count is the number of rows deleted. Note that the number may be less than the number of rows that matched the condition when deletes were suppressed by a BEFORE DELETE trigger. If count is 0, no rows were deleted by the query (this is not considered an error).

If the DELETE command contains a RETURNING clause, the result will be similar to that of a SELECT statement containing the columns and values defined in the RETURNING list, computed over the row(s) deleted by the command.

Notes

PostgreSQL lets you reference columns of other tables in the WHERE condition by specifying the other tables in the USING clause. For example, to delete all films produced by a given producer, one can do:

```
DELETE FROM films USING producers
    WHERE producer_id = producers.id AND producers.name = 'foo';
```
What is essentially happening here is a join between films and producers, with all successfully joined films rows being marked for deletion. This syntax is not standard. A more standard way to do it is:

```
DELETE FROM films
    WHERE producer_id IN (SELECT id FROM producers WHERE name = 'foo');
```

In some cases the join style is easier to write or faster to execute than the sub-select style.

**Examples**

Delete all films but musicals:

```
DELETE FROM films WHERE kind <> 'Musical';
```

Clear the table films:

```
DELETE FROM films;
```

Delete completed tasks, returning full details of the deleted rows:

```
DELETE FROM tasks WHERE status = 'DONE' RETURNING *;
```

Delete the row of tasks on which the cursor c_tasks is currently positioned:

```
DELETE FROM tasks WHERE CURRENT OF c_tasks;
```

**Compatibility**

This command conforms to the SQL standard, except that the using and returning clauses are PostgreSQL extensions, as is the ability to use with with delete.
**DISCARD**

### Name

DISCARD — discard session state

### Synopsis

DISCARD ( ALL | PLANS | SEQUENCES | TEMPORARY | TEMP )

### Description

DISCARD releases internal resources associated with a database session. This command is useful for partially or fully resetting the session’s state. There are several subcommands to release different types of resources; the DISCARD ALL variant subsumes all the others, and also resets additional state.

### Parameters

**PLANS**

Releases all cached query plans, forcing re-planning to occur the next time the associated prepared statement is used.

**SEQUENCES**

Discards all cached sequence-related state, including currval()/lastval() information and any preallocated sequence values that have not yet been returned by nextval(). (See CREATE SEQUENCE for a description of preallocated sequence values.)

**TEMPORARY or TEMP**

Drops all temporary tables created in the current session.

**ALL**

Releases all temporary resources associated with the current session and resets the session to its initial state. Currently, this has the same effect as executing the following sequence of statements:

```sql
SET SESSION AUTHORIZATION DEFAULT;
RESET ALL;
DEALLOCATE ALL;
CLOSE ALL;
UNLISTEN *;
SELECT pg_advisory_unlock_all();
DISCARD PLANS;
DISCARD SEQUENCES;
DISCARD TEMP;
```

### Notes

DISCARD ALL cannot be executed inside a transaction block.
Compatibility

DISCARD is a PostgreSQL extension.
**DO**

**Name**

DO — execute an anonymous code block

**Synopsis**

DO [ LANGUAGE lang_name ] code

**Description**

DO executes an anonymous code block, or in other words a transient anonymous function in a procedural language.

The code block is treated as though it were the body of a function with no parameters, returning `void`. It is parsed and executed a single time.

The optional `LANGUAGE` clause can be written either before or after the code block.

**Parameters**

- **code**
  
  The procedural language code to be executed. This must be specified as a string literal, just as in `CREATE FUNCTION`. Use of a dollar-quoted literal is recommended.

- **lang_name**
  
  The name of the procedural language the code is written in. If omitted, the default is `plpgsql`.

**Notes**

The procedural language to be used must already have been installed into the current database by means of `CREATE LANGUAGE`. `plpgsql` is installed by default, but other languages are not.

The user must have `USAGE` privilege for the procedural language, or must be a superuser if the language is untrusted. This is the same privilege requirement as for creating a function in the language.

**Examples**

Grant all privileges on all views in schema **public** to role **webuser**:

```sql
DO $$
DECLARE r record;
BEGIN
  FOR r IN SELECT table_schema, table_name FROM information_schema.tables
  WHERE table_type = 'VIEW' AND table_schema = 'public'
  LOOP
    EXECUTE 'GRANT ALL ON ' || quote_ident(r.table_schema) || '.' || quote_ident(r.t.
  END LOOP;
$$
```
DO

END\$$;

Compatibility

There is no `DO` statement in the SQL standard.

See Also

CREATE LANGUAGE
DROP ACCESS METHOD

Name
DROP ACCESS METHOD — remove an access method

Synopsis
DROP ACCESS METHOD [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description
DROP ACCESS METHOD removes an existing access method. Only superusers can drop access methods.

Parameters
IF EXISTS
Do not throw an error if the access method does not exist. A notice is issued in this case.

name
The name of an existing access method.

CASCADE
Automatically drop objects that depend on the access method (such as operator classes, operator families, and indexes), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the access method if any objects depend on it. This is the default.

Examples
Drop the access method heptree:

DROP ACCESS METHOD heptree;

Compatibility
DROP ACCESS METHOD is a PostgreSQL extension.

See Also
CREATE ACCESS METHOD
DROP AGGREGATE

Name

DROP AGGREGATE — remove an aggregate function

Synopsis

DROP AGGREGATE [ IF EXISTS ] name ( aggregate_signature ) [ CASCADE | RESTRICT ]

where aggregate_signature is:

* | [ argmode ] [ argname ] argtype [ , ... ] |
[ [ argmode ] [ argname ] argtype [ , ... ] ] ORDER BY [ argmode ] [ argname ] argtype [ , ... ]

Description

DROP AGGREGATE removes an existing aggregate function. To execute this command the current user must be the owner of the aggregate function.

Parameters

IF EXISTS

Do not throw an error if the aggregate does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing aggregate function.

argmode

The mode of an argument: IN or VARIADIC. If omitted, the default is IN.

argname

The name of an argument. Note that DROP AGGREGATE does not actually pay any attention to argument names, since only the argument data types are needed to determine the aggregate function’s identity.

argtype

An input data type on which the aggregate function operates. To reference a zero-argument aggregate function, write * in place of the list of argument specifications. To reference an ordered-set aggregate function, write ORDER BY between the direct and aggregated argument specifications.

CASCADE

Automatically drop objects that depend on the aggregate function (such as views using it), and in turn all objects that depend on those objects (see Section 5.13).
DROP AGGREGATE

Restrict
Refuse to drop the aggregate function if any objects depend on it. This is the default.

Notes
Alternative syntaxes for referencing ordered-set aggregates are described under ALTER AGGREGATE.

Examples
To remove the aggregate function myavg for type integer:

DROP AGGREGATE myavg(integer);

To remove the hypothetical-set aggregate function myrank, which takes an arbitrary list of ordering columns and a matching list of direct arguments:

DROP AGGREGATE myrank(VARIADIC "any" ORDER BY VARIADIC "any");

Compatibility
There is no DROP AGGREGATE statement in the SQL standard.

See Also
ALTER AGGREGATE, CREATE AGGREGATE
DROP CAST

Name

DROP CAST — remove a cast

Synopsis

DROP CAST [ IF EXISTS ] (source_type AS target_type) [ CASCADE | RESTRICT ]

Description

DROP CAST removes a previously defined cast.
To be able to drop a cast, you must own the source or the target data type. These are the same privileges that are required to create a cast.

Parameters

IF EXISTS
Do not throw an error if the cast does not exist. A notice is issued in this case.

source_type
The name of the source data type of the cast.

target_type
The name of the target data type of the cast.

CASCADE
RESTRICT
These key words do not have any effect, since there are no dependencies on casts.

Examples

To drop the cast from type text to type int:

DROP CAST (text AS int);

Compatibility

The DROP CAST command conforms to the SQL standard.
See Also

CREATE CAST
DROP COLLATION

Name
DROP COLLATION — remove a collation

Synopsis
DROP COLLATION [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description
DROP COLLATION removes a previously defined collation. To be able to drop a collation, you must own the collation.

Parameters
IF EXISTS
Do not throw an error if the collation does not exist. A notice is issued in this case.

name
The name of the collation. The collation name can be schema-qualified.

CASCADE
Automatically drop objects that depend on the collation, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the collation if any objects depend on it. This is the default.

Examples
To drop the collation named german:

DROP COLLATION german;

Compatibility
The DROP COLLATION command conforms to the SQL standard, apart from the IF EXISTS option, which is a PostgreSQL extension.
See Also

ALTER COLLATION, CREATE COLLATION
DROP CONVERSION

Name
DROP CONVERSION — remove a conversion

Synopsis
DROP CONVERSION [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description
DROP CONVERSION removes a previously defined conversion. To be able to drop a conversion, you must own the conversion.

Parameters
 IF EXISTS
 Do not throw an error if the conversion does not exist. A notice is issued in this case.

 name
 The name of the conversion. The conversion name can be schema-qualified.

 CASCADE
 RESTRICT
 These key words do not have any effect, since there are no dependencies on conversions.

Examples
To drop the conversion named myname:

DROP CONVERSION myname;

Compatibility
There is no DROP CONVERSION statement in the SQL standard, but a DROP TRANSLATION statement that goes along with the CREATE TRANSLATION statement that is similar to the CREATE CONVERSION statement in PostgreSQL.

See Also
ALTER CONVERSION, CREATE CONVERSION
DROP DATABASE

Name
DROP DATABASE — remove a database

Synopsis
DROP DATABASE [ IF EXISTS ] name

Description
DROP DATABASE drops a database. It removes the catalog entries for the database and deletes the directory containing the data. It can only be executed by the database owner. Also, it cannot be executed while you or anyone else are connected to the target database. (Connect to postgres or any other database to issue this command.)

DROP DATABASE cannot be undone. Use it with care!

Parameters

IF EXISTS

Do not throw an error if the database does not exist. A notice is issued in this case.

name

The name of the database to remove.

Notes

DROP DATABASE cannot be executed inside a transaction block.

This command cannot be executed while connected to the target database. Thus, it might be more convenient to use the program dropdb instead, which is a wrapper around this command.

Compatibility

There is no DROP DATABASE statement in the SQL standard.

See Also

CREATE DATABASE
DROP DOMAIN

Name
DROP DOMAIN — remove a domain

Synopsis
DROP DOMAIN [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description
DROP DOMAIN removes a domain. Only the owner of a domain can remove it.

Parameters

IF EXISTS
Do not throw an error if the domain does not exist. A notice is issued in this case.

name
The name (optionally schema-qualified) of an existing domain.

CASCADE
Automatically drop objects that depend on the domain (such as table columns), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the domain if any objects depend on it. This is the default.

Examples
To remove the domain box:

DROP DOMAIN box;

Compatibility
This command conforms to the SQL standard, except for the IF EXISTS option, which is a PostgreSQL extension.

See Also
CREATE DOMAIN, ALTER DOMAIN
DROP EVENT TRIGGER

Name
DROP EVENT TRIGGER — remove an event trigger

Synopsis
DROP EVENT TRIGGER [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description
DROP EVENT TRIGGER removes an existing event trigger. To execute this command, the current user must be the owner of the event trigger.

Parameters

IF EXISTS
Do not throw an error if the event trigger does not exist. A notice is issued in this case.

name
The name of the event trigger to remove.

CASCADE
Automatically drop objects that depend on the trigger, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the trigger if any objects depend on it. This is the default.

Examples
Destroy the trigger snitch:

DROP EVENT TRIGGER snitch;

Compatibility
There is no DROP EVENT TRIGGER statement in the SQL standard.

See Also
CREATE EVENT TRIGGER, ALTER EVENT TRIGGER
DROP EXTENSION

Name
DROP EXTENSION — remove an extension

Synopsis
DROP EXTENSION [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description
DROP EXTENSION removes extensions from the database. Dropping an extension causes its component objects to be dropped as well.
You must own the extension to use DROP EXTENSION.

Parameters

IF EXISTS
Do not throw an error if the extension does not exist. A notice is issued in this case.

name
The name of an installed extension.

CASCADE
Automatically drop objects that depend on the extension, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the extension if any objects depend on it (other than its own member objects and other extensions listed in the same DROP command). This is the default.

Examples
To remove the extension hstore from the current database:

DROP EXTENSION hstore;

This command will fail if any of hstore’s objects are in use in the database, for example if any tables have columns of the hstore type. Add the CASCADE option to forcibly remove those dependent objects as well.

Compatibility
DROP EXTENSION is a PostgreSQL extension.
See Also

CREATE EXTENSION, ALTER EXTENSION
DROP FOREIGN DATA WRAPPER

Name
DROP FOREIGN DATA WRAPPER — remove a foreign-data wrapper

Synopsis
DROP FOREIGN DATA WRAPPER [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description
DROP FOREIGN DATA WRAPPER removes an existing foreign-data wrapper. To execute this command, the current user must be the owner of the foreign-data wrapper.

Parameters
IF EXISTS
Do not throw an error if the foreign-data wrapper does not exist. A notice is issued in this case.

name
The name of an existing foreign-data wrapper.

CASCADE
Automatically drop objects that depend on the foreign-data wrapper (such as foreign tables and servers), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the foreign-data wrapper if any objects depend on it. This is the default.

Examples
Drop the foreign-data wrapper dbi:

DROP FOREIGN DATA WRAPPER dbi;

Compatibility
DROP FOREIGN DATA WRAPPER conforms to ISO/IEC 9075-9 (SQL/MED). The IF EXISTS clause is a PostgreSQL extension.
See Also

CREATE FOREIGN DATA WRAPPER, ALTER FOREIGN DATA WRAPPER
DROP FOREIGN TABLE

**Name**

DROP FOREIGN TABLE — remove a foreign table

**Synopsis**

DROP FOREIGN TABLE [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

**Description**

DROP FOREIGN TABLE removes a foreign table. Only the owner of a foreign table can remove it.

**Parameters**

**IF EXISTS**

Do not throw an error if the foreign table does not exist. A notice is issued in this case.

**name**

The name (optionally schema-qualified) of the foreign table to drop.

**CASCADE**

Automatically drop objects that depend on the foreign table (such as views), and in turn all objects that depend on those objects (see Section 5.13).

**RESTRICT**

Refuse to drop the foreign table if any objects depend on it. This is the default.

**Examples**

To destroy two foreign tables, films and distributors:

```
DROP FOREIGN TABLE films, distributors;
```

**Compatibility**

This command conforms to the ISO/IEC 9075-9 (SQL/MED), except that the standard only allows one foreign table to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.
See Also

ALTER FOREIGN TABLE, CREATE FOREIGN TABLE
DROP FUNCTION

Name

DROP FUNCTION — remove a function

Synopsis

DROP FUNCTION [ IF EXISTS ] name ( [ [ argmode ] [ argname ] argtype [, ...] ] )
[ CASCADE | RESTRICT ]

Description

DROP FUNCTION removes the definition of an existing function. To execute this command the user must be the owner of the function. The argument types to the function must be specified, since several different functions can exist with the same name and different argument lists.

Parameters

IF EXISTS

Do not throw an error if the function does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing function.

argmode

The mode of an argument: IN, OUT, INOUT, or VARIADIC. If omitted, the default is IN. Note that DROP FUNCTION does not actually pay any attention to OUT arguments, since only the input arguments are needed to determine the function’s identity. So it is sufficient to list the IN, INOUT, and VARIADIC arguments.

argname

The name of an argument. Note that DROP FUNCTION does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

argtype

The data type(s) of the function’s arguments (optionally schema-qualified), if any.

CASCADE

Automatically drop objects that depend on the function (such as operators or triggers), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the function if any objects depend on it. This is the default.
Examples

This command removes the square root function:

\texttt{DROP FUNCTION sqrt(integer);}

Compatibility

A \texttt{DROP FUNCTION} statement is defined in the SQL standard, but it is not compatible with this command.

See Also

\texttt{CREATE FUNCTION, ALTER FUNCTION}
DROP GROUP

Name
DROP GROUP — remove a database role

Synopsis
DROP GROUP [ IF EXISTS ] name [, ...]

Description
DROP GROUP is now an alias for DROP ROLE.

Compatibility
There is no DROP GROUP statement in the SQL standard.

See Also
DROP ROLE
**DROP INDEX**

**Name**

DROP INDEX — remove an index

**Synopsis**

\[
\text{DROP INDEX [ CONCURRENTLY ] [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]}
\]

**Description**

DROP INDEX drops an existing index from the database system. To execute this command you must be the owner of the index.

**Parameters**

**CONCURRENTLY**

Drop the index without locking out concurrent selects, inserts, updates, and deletes on the index’s table. A normal DROP INDEX acquires exclusive lock on the table, blocking other accesses until the index drop can be completed. With this option, the command instead waits until conflicting transactions have completed.

There are several caveats to be aware of when using this option. Only one index name can be specified, and the CASCADE option is not supported. (Thus, an index that supports a UNIQUE or PRIMARY KEY constraint cannot be dropped this way.) Also, regular DROP INDEX commands can be performed within a transaction block, but DROP INDEX CONCURRENTLY cannot.

For temporary tables, DROP INDEX is always non-concurrent, as no other session can access them, and non-concurrent index drop is cheaper.

**IF EXISTS**

Do not throw an error if the index does not exist. A notice is issued in this case.

**name**

The name (optionally schema-qualified) of an index to remove.

**CASCADE**

Automatically drop objects that depend on the index, and in turn all objects that depend on those objects (see Section 5.13).

**RESTRICT**

Refuse to drop the index if any objects depend on it. This is the default.
Examples

This command will remove the index \texttt{title_idx}:

\texttt{DROP INDEX title_idx;} 

Compatibility

\texttt{DROP INDEX} is a PostgreSQL language extension. There are no provisions for indexes in the SQL standard.

See Also

\texttt{CREATE INDEX}
DROP LANGUAGE

Name

DROP LANGUAGE — remove a procedural language

Synopsis

DROP [ PROCEDURAL ] LANGUAGE [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP LANGUAGE removes the definition of a previously registered procedural language. You must be a superuser or the owner of the language to use DROP LANGUAGE.

Note: As of PostgreSQL 9.1, most procedural languages have been made into “extensions”, and should therefore be removed with DROP EXTENSION not DROP LANGUAGE.

Parameters

IF EXISTS

Do not throw an error if the language does not exist. A notice is issued in this case.

name

The name of an existing procedural language. For backward compatibility, the name can be enclosed by single quotes.

CASCADE

Automatically drop objects that depend on the language (such as functions in the language), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the language if any objects depend on it. This is the default.

Examples

This command removes the procedural language plsample:

DROP LANGUAGE plsample;
Compatibility

There is no DROP LANGUAGE statement in the SQL standard.

See Also

ALTER LANGUAGE, CREATE LANGUAGE, droplang
DROP MATERIALIZED VIEW

Name
DROP MATERIALIZED VIEW — remove a materialized view

Synopsis
DROP MATERIALIZED VIEW [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description
DROP MATERIALIZED VIEW drops an existing materialized view. To execute this command you must be the owner of the materialized view.

Parameters
IF EXISTS
Do not throw an error if the materialized view does not exist. A notice is issued in this case.

name
The name (optionally schema-qualified) of the materialized view to remove.

CASCADE
Automatically drop objects that depend on the materialized view (such as other materialized views, or regular views), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the materialized view if any objects depend on it. This is the default.

Examples
This command will remove the materialized view called order_summary:

DROP MATERIALIZED VIEW order_summary;

Compatibility
DROP MATERIALIZED VIEW is a PostgreSQL extension.
See Also

CREATE MATERIALIZED VIEW, ALTER MATERIALIZED VIEW, REFRESH MATERIALIZED VIEW
DROP OPERATOR

Name

DROP OPERATOR — remove an operator

Synopsis

DROP OPERATOR [ IF EXISTS ] name ( { left_type | NONE } , { right_type | NONE } ) [ CASCADE | RESTRICT ]

Description

DROP OPERATOR drops an existing operator from the database system. To execute this command you must be the owner of the operator.

Parameters

IF EXISTS

Do not throw an error if the operator does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing operator.

left_type

The data type of the operator’s left operand; write NONE if the operator has no left operand.

right_type

The data type of the operator’s right operand; write NONE if the operator has no right operand.

CASCADE

Automatically drop objects that depend on the operator (such as views using it), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the operator if any objects depend on it. This is the default.

Examples

Remove the power operator $a^b$ for type integer:

DROP OPERATOR ^ (integer, integer);

Remove the left unary bitwise complement operator ~b for type bit:

DROP OPERATOR ~ (none, bit);
Remove the right unary factorial operator $x!$ for type `bigint`:

```sql
DROP OPERATOR ! (bigint, none);
```

**Compatibility**

There is no `DROP OPERATOR` statement in the SQL standard.

**See Also**

CREATE OPERATOR, ALTER OPERATOR
DROP OPERATOR CLASS

Name

DROP OPERATOR CLASS — remove an operator class

Synopsis

DROP OPERATOR CLASS [ IF EXISTS ] name USING index_method [ CASCADE | RESTRICT ]

Description

DROP OPERATOR CLASS drops an existing operator class. To execute this command you must be the owner of the operator class.

DROP OPERATOR CLASS does not drop any of the operators or functions referenced by the class. If there are any indexes depending on the operator class, you will need to specify CASCADE for the drop to complete.

Parameters

IF EXISTS

Do not throw an error if the operator class does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing operator class.

index_method

The name of the index access method the operator class is for.

CASCADE

Automatically drop objects that depend on the operator class (such as indexes), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the operator class if any objects depend on it. This is the default.

Notes

DROP OPERATOR CLASS will not drop the operator family containing the class, even if there is nothing else left in the family (in particular, in the case where the family was implicitly created by CREATE OPERATOR CLASS). An empty operator family is harmless, but for the sake of tidiness you might wish to remove the family with DROP OPERATOR FAMILY; or perhaps better, use DROP OPERATOR FAMILY in the first place.
Examples

Remove the B-tree operator class `widget_ops`:

```
DROP OPERATOR CLASS widget_ops USING btree;
```

This command will not succeed if there are any existing indexes that use the operator class. Add `CASCADE` to drop such indexes along with the operator class.

Compatibility

There is no `DROP OPERATOR CLASS` statement in the SQL standard.

See Also

`ALTER OPERATOR CLASS`, `CREATE OPERATOR CLASS`, `DROP OPERATOR FAMILY`
DROP OPERATOR FAMILY

Name
DROP OPERATOR FAMILY — remove an operator family

Synopsis
DROP OPERATOR FAMILY [ IF EXISTS ] name USING index_method [ CASCADE | RESTRICT ]

Description
DROP OPERATOR FAMILY drops an existing operator family. To execute this command you must be the owner of the operator family.

DROP OPERATOR FAMILY includes dropping any operator classes contained in the family, but it does not drop any of the operators or functions referenced by the family. If there are any indexes depending on operator classes within the family, you will need to specify CASCADE for the drop to complete.

Parameters
IF EXISTS
Do not throw an error if the operator family does not exist. A notice is issued in this case.

name
The name (optionally schema-qualified) of an existing operator family.

index_method
The name of the index access method the operator family is for.

CASCADE
Automatically drop objects that depend on the operator family, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the operator family if any objects depend on it. This is the default.

Examples
Remove the B-tree operator family float_ops:

DROP OPERATOR FAMILY float_ops USING btree;

This command will not succeed if there are any existing indexes that use operator classes within the family. Add CASCADE to drop such indexes along with the operator family.
Compatibility

There is no `DROP OPERATOR FAMILY` statement in the SQL standard.

See Also

`ALTER OPERATOR FAMILY`, `CREATE OPERATOR FAMILY`, `ALTER OPERATOR CLASS`, `CREATE OPERATOR CLASS`, `DROP OPERATOR CLASS`
DROP OWNED

Name
DROP OWNED — remove database objects owned by a database role

Synopsis
DROP OWNED BY { name | CURRENT_USER | SESSION_USER } [, ...) [ CASCADE | RESTRICT ]

Description
DROP OWNED drops all the objects within the current database that are owned by one of the specified roles. Any privileges granted to the given roles on objects in the current database or on shared objects (databases, tablespaces) will also be revoked.

Parameters
name
The name of a role whose objects will be dropped, and whose privileges will be revoked.

CASCADE
Automatically drop objects that depend on the affected objects, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the objects owned by a role if any other database objects depend on one of the affected objects. This is the default.

Notes
DROP OWNED is often used to prepare for the removal of one or more roles. Because DROP OWNED only affects the objects in the current database, it is usually necessary to execute this command in each database that contains objects owned by a role that is to be removed.

Using the CASCADE option might make the command recurse to objects owned by other users.

The REASSIGN OWNED command is an alternative that reassigns the ownership of all the database objects owned by one or more roles. However, REASSIGN OWNED does not deal with privileges for other objects.

Databases and tablespaces owned by the role(s) will not be removed.

See Section 21.4 for more discussion.
Compatibility

The `DROP OWNED` command is a PostgreSQL extension.

See Also

`REASSIGN OWNED`, `DROP ROLE`
DROP POLICY

Name
DROP POLICY — remove a row level security policy from a table

Synopsis
DROP POLICY [ IF EXISTS ] name ON table_name [ CASCADE | RESTRICT ]

Description
DROP POLICY removes the specified policy from the table. Note that if the last policy is removed for a table and the table still has row level security enabled via ALTER TABLE, then the default-deny policy will be used. ALTER TABLE ... DISABLE ROW LEVEL SECURITY can be used to disable row level security for a table, whether policies for the table exist or not.

Parameters
IF EXISTS
Do not throw an error if the policy does not exist. A notice is issued in this case.

name
The name of the policy to drop.

table_name
The name (optionally schema-qualified) of the table that the policy is on.

CASCADE
RESTRICT
These key words do not have any effect, since there are no dependencies on policies.

Examples
To drop the policy called p1 on the table named my_table:

DROP POLICY p1 ON my_table;

Compatibility
DROP POLICY is a PostgreSQL extension.
See Also

CREATE POLICY, ALTER POLICY
DROP ROLE

Name
DROP ROLE — remove a database role

Synopsis
DROP ROLE [ IF EXISTS ] name [, ...]

Description
DROP ROLE removes the specified role(s). To drop a superuser role, you must be a superuser yourself; to drop non-superuser roles, you must have CREATEROLE privilege. A role cannot be removed if it is still referenced in any database of the cluster; an error will be raised if so. Before dropping the role, you must drop all the objects it owns (or reassign their ownership) and revoke any privileges the role has been granted on other objects. The REASSIGN OWNED and DROP OWNED commands can be useful for this purpose; see Section 21.4 for more discussion. However, it is not necessary to remove role memberships involving the role; DROP ROLE automatically revokes any memberships of the target role in other roles, and of other roles in the target role. The other roles are not dropped nor otherwise affected.

Parameters
IF EXISTS
Do not throw an error if the role does not exist. A notice is issued in this case.

name
The name of the role to remove.

Notes
PostgreSQL includes a program dropuser that has the same functionality as this command (in fact, it calls this command) but can be run from the command shell.

Examples
To drop a role:
DROP ROLE jonathan;
Compatibility
The SQL standard defines `DROP ROLE`, but it allows only one role to be dropped at a time, and it specifies different privilege requirements than PostgreSQL uses.

See Also
CREATE ROLE, ALTER ROLE, SET ROLE
DROP RULE

**Name**

DROP RULE — remove a rewrite rule

**Synopsis**

DROP RULE [ IF EXISTS ] name ON table_name [ CASCADE | RESTRICT ]

**Description**

DROP RULE drops a rewrite rule.

**Parameters**

**IF EXISTS**

Do not throw an error if the rule does not exist. A notice is issued in this case.

**name**

The name of the rule to drop.

**table_name**

The name (optionally schema-qualified) of the table or view that the rule applies to.

**CASCADE**

Automatically drop objects that depend on the rule, and in turn all objects that depend on those objects (see Section 5.13).

**RESTRICT**

Refuse to drop the rule if any objects depend on it. This is the default.

**Examples**

To drop the rewrite rule `newrule`:

DROP RULE newrule ON mytable;

**Compatibility**

DROP RULE is a PostgreSQL language extension, as is the entire query rewrite system.
See Also

CREATE RULE, ALTER RULE
DROP SCHEMA

Name
DROP SCHEMA — remove a schema

Synopsis
DROP SCHEMA [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description
DROP SCHEMA removes schemas from the database.
A schema can only be dropped by its owner or a superuser. Note that the owner can drop the schema (and thereby all contained objects) even if they do not own some of the objects within the schema.

Parameters

IF EXISTS
Do not throw an error if the schema does not exist. A notice is issued in this case.

name
The name of a schema.

CASCADE
Automatically drop objects (tables, functions, etc.) that are contained in the schema, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the schema if it contains any objects. This is the default.

Notes
Using the CASCADE option might make the command remove objects in other schemas besides the one(s) named.

Examples
To remove schema mystuff from the database, along with everything it contains:
DROP SCHEMA mystuff CASCADE;
Compatibility

DROP SCHEMA is fully conforming with the SQL standard, except that the standard only allows one schema to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also

ALTER SCHEMA, CREATE SCHEMA
DROP SEQUENCE

Name

DROP SEQUENCE — remove a sequence

Synopsis

DROP SEQUENCE [ IF EXISTS ] name [, ... ] [ CASCADE | RESTRICT ]

Description

DROP SEQUENCE removes sequence number generators. A sequence can only be dropped by its owner or a superuser.

Parameters

IF EXISTS

Do not throw an error if the sequence does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of a sequence.

CASCADE

Automatically drop objects that depend on the sequence, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the sequence if any objects depend on it. This is the default.

Examples

To remove the sequence serial:

DROP SEQUENCE serial;

Compatibility

DROP SEQUENCE conforms to the SQL standard, except that the standard only allows one sequence to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.
See Also

CREATE SEQUENCE, ALTER SEQUENCE
**DROP SERVER**

**Name**

DROP SERVER — remove a foreign server descriptor

**Synopsis**

DROP SERVER [ IF EXISTS ] name [ CASCADE | RESTRICT ]

**Description**

DROP SERVER removes an existing foreign server descriptor. To execute this command, the current user must be the owner of the server.

**Parameters**

- **IF EXISTS**
  Do not throw an error if the server does not exist. A notice is issued in this case.

- **name**
  The name of an existing server.

- **CASCADE**
  Automatically drop objects that depend on the server (such as user mappings), and in turn all objects that depend on those objects (see Section 5.13).

- **RESTRICT**
  Refuse to drop the server if any objects depend on it. This is the default.

**Examples**

Drop a server foo if it exists:

```
DROP SERVER IF EXISTS foo;
```

**Compatibility**

DROP SERVER conforms to ISO/IEC 9075-9 (SQL/MED). The IF EXISTS clause is a PostgreSQL extension.
See Also

CREATE SERVER, ALTER SERVER
DROP TABLE

Name

DROP TABLE — remove a table

Synopsis

DROP TABLE [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description

DROP TABLE removes tables from the database. Only the table owner, the schema owner, and superuser can drop a table. To empty a table of rows without destroying the table, use DELETE or TRUNCATE.

DROP TABLE always removes any indexes, rules, triggers, and constraints that exist for the target table. However, to drop a table that is referenced by a view or a foreign-key constraint of another table, CASCADE must be specified. (CASCADE will remove a dependent view entirely, but in the foreign-key case it will only remove the foreign-key constraint, not the other table entirely.)

Parameters

IF EXISTS

Do not throw an error if the table does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of the table to drop.

CASCADE

Automatically drop objects that depend on the table (such as views), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the table if any objects depend on it. This is the default.

Examples

To destroy two tables, films and distributors:

DROP TABLE films, distributors;
Compatibility

This command conforms to the SQL standard, except that the standard only allows one table to be dropped per command, and apart from the `IF EXISTS` option, which is a PostgreSQL extension.

See Also

ALTER TABLE, CREATE TABLE
DROP TABLESPACE

Name
DROP TABLESPACE — remove a tablespace

Synopsis
DROP TABLESPACE [ IF EXISTS ] name

Description
DROP TABLESPACE removes a tablespace from the system.
A tablespace can only be dropped by its owner or a superuser. The tablespace must be empty of all
database objects before it can be dropped. It is possible that objects in other databases might still
reside in the tablespace even if no objects in the current database are using the tablespace. Also, if the
tablespace is listed in the temp_tablespaces setting of any active session, the DROP might fail due to
temporary files residing in the tablespace.

Parameters
IF EXISTS
Do not throw an error if the tablespace does not exist. A notice is issued in this case.
name
The name of a tablespace.

Notes
DROP TABLESPACE cannot be executed inside a transaction block.

Examples
To remove tablespace mystuff from the system:
DROP TABLESPACE mystuff;

Compatibility
DROP TABLESPACE is a PostgreSQL extension.
DROP TABLESPACE

See Also
CREATE TABLESPACE, ALTER TABLESPACE
DROP TEXT SEARCH CONFIGURATION

Name

DROP TEXT SEARCH CONFIGURATION — remove a text search configuration

Synopsis

DROP TEXT SEARCH CONFIGURATION [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP TEXT SEARCH CONFIGURATION drops an existing text search configuration. To execute this command you must be the owner of the configuration.

Parameters

IF EXISTS

Do not throw an error if the text search configuration does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing text search configuration.

CASCADE

Automatically drop objects that depend on the text search configuration, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the text search configuration if any objects depend on it. This is the default.

Examples

Remove the text search configuration my_english:

DROP TEXT SEARCH CONFIGURATION my_english;

This command will not succeed if there are any existing indexes that reference the configuration in to_tsvector calls. Add CASCADE to drop such indexes along with the text search configuration.

Compatibility

There is no DROP TEXT SEARCH CONFIGURATION statement in the SQL standard.
See Also

ALTER TEXT SEARCH CONFIGURATION, CREATE TEXT SEARCH CONFIGURATION
DROP TEXT SEARCH DICTIONARY

Name
DROP TEXT SEARCH DICTIONARY — remove a text search dictionary

Synopsis
DROP TEXT SEARCH DICTIONARY [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description
DROP TEXT SEARCH DICTIONARY drops an existing text search dictionary. To execute this command you must be the owner of the dictionary.

Parameters
IF EXISTS
Do not throw an error if the text search dictionary does not exist. A notice is issued in this case.
name
The name (optionally schema-qualified) of an existing text search dictionary.
CASCADE
Automatically drop objects that depend on the text search dictionary, and in turn all objects that depend on those objects (see Section 5.13).
RESTRICT
Refuse to drop the text search dictionary if any objects depend on it. This is the default.

Examples
Remove the text search dictionary english:
DROP TEXT SEARCH DICTIONARY english;

This command will not succeed if there are any existing text search configurations that use the dictionary. Add CASCADE to drop such configurations along with the dictionary.

Compatibility
There is no DROP TEXT SEARCH DICTIONARY statement in the SQL standard.
See Also

ALTER TEXT SEARCH DICTIONARY, CREATE TEXT SEARCH DICTIONARY
DROP TEXT SEARCH PARSER

Name

DROP TEXT SEARCH PARSER — remove a text search parser

Synopsis

DROP TEXT SEARCH PARSER [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description

DROP TEXT SEARCH PARSER drops an existing text search parser. You must be a superuser to use this command.

Parameters

IF EXISTS

Do not throw an error if the text search parser does not exist. A notice is issued in this case.

name

The name (optionally schema-qualified) of an existing text search parser.

CASCADE

Automatically drop objects that depend on the text search parser, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the text search parser if any objects depend on it. This is the default.

Examples

Remove the text search parser my_parser:

DROP TEXT SEARCH PARSER my_parser;

This command will not succeed if there are any existing text search configurations that use the parser. Add CASCADE to drop such configurations along with the parser.

Compatibility

There is no DROP TEXT SEARCH PARSER statement in the SQL standard.
See Also

ALTER TEXT SEARCH PARSER, CREATE TEXT SEARCH PARSER
DROP TEXT SEARCH TEMPLATE

Name
DROP TEXT SEARCH TEMPLATE — remove a text search template

Synopsis
DROP TEXT SEARCH TEMPLATE [ IF EXISTS ] name [ CASCADE | RESTRICT ]

Description
DROP TEXT SEARCH TEMPLATE drops an existing text search template. You must be a superuser to use this command.

Parameters
IF EXISTS
Do not throw an error if the text search template does not exist. A notice is issued in this case.

name
The name (optionally schema-qualified) of an existing text search template.

CASCADE
Automatically drop objects that depend on the text search template, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the text search template if any objects depend on it. This is the default.

Examples
Remove the text search template thesaurus:

DROP TEXT SEARCH TEMPLATE thesaurus;

This command will not succeed if there are any existing text search dictionaries that use the template. Add CASCADE to drop such dictionaries along with the template.

Compatibility
There is no DROP TEXT SEARCH TEMPLATE statement in the SQL standard.
See Also

ALTER TEXT SEARCH TEMPLATE, CREATE TEXT SEARCH TEMPLATE
DROP TRANSFORM

Name

DROP TRANSFORM — remove a transform

Synopsis

DROP TRANSFORM [ IF EXISTS ] FOR type_name LANGUAGE lang_name [ CASCADE | RESTRICT ]

Description

DROP TRANSFORM removes a previously defined transform.
To be able to drop a transform, you must own the type and the language. These are the same privileges that are required to create a transform.

Parameters

IF EXISTS

Do not throw an error if the transform does not exist. A notice is issued in this case.

type_name

The name of the data type of the transform.

lang_name

The name of the language of the transform.

CASCADE

Automatically drop objects that depend on the transform, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT

Refuse to drop the transform if any objects depend on it. This is the default.

Examples

To drop the transform for type hstore and language plpythonu:

DROP TRANSFORM FOR hstore LANGUAGE plpythonu;

Compatibility

This form of DROP TRANSFORM is a PostgreSQL extension. See CREATE TRANSFORM for details.
See Also

CREATE TRANSFORM
DROP TRIGGER

Name
DROP TRIGGER — remove a trigger

Synopsis
DROP TRIGGER [ IF EXISTS ] name ON table_name [ CASCADE | RESTRICT ]

Description
DROP TRIGGER removes an existing trigger definition. To execute this command, the current user must be the owner of the table for which the trigger is defined.

Parameters

IF EXISTS
Do not throw an error if the trigger does not exist. A notice is issued in this case.

name
The name of the trigger to remove.

table_name
The name (optionally schema-qualified) of the table for which the trigger is defined.

CASCADE
Automatically drop objects that depend on the trigger, and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the trigger if any objects depend on it. This is the default.

Examples
Destroy the trigger if_dist_exists on the table films:

DROP TRIGGER if_dist_exists ON films;

Compatibility
The DROP TRIGGER statement in PostgreSQL is incompatible with the SQL standard. In the SQL standard, trigger names are not local to tables, so the command is simply DROP TRIGGER name.
See Also

CREATE TRIGGER
DROP TYPE

Name
DROP TYPE — remove a data type

Synopsis
DROP TYPE [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description
DROP TYPE removes a user-defined data type. Only the owner of a type can remove it.

Parameters
IF EXISTS
Do not throw an error if the type does not exist. A notice is issued in this case.

name
The name (optionally schema-qualified) of the data type to remove.

CASCADE
Automatically drop objects that depend on the type (such as table columns, functions, and operators), and in turn all objects that depend on those objects (see Section 5.13).

RESTRICT
Refuse to drop the type if any objects depend on it. This is the default.

Examples
To remove the data type box:

DROP TYPE box;

Compatibility
This command is similar to the corresponding command in the SQL standard, apart from the IF EXISTS option, which is a PostgreSQL extension. But note that much of the CREATE TYPE command and the data type extension mechanisms in PostgreSQL differ from the SQL standard.
See Also

ALTER TYPE, CREATE TYPE
DROP USER

Name
DROP USER — remove a database role

Synopsis
DROP USER [ IF EXISTS ] name [, ...]

Description
DROP USER is simply an alternate spelling of DROP ROLE.

Compatibility
The DROP USER statement is a PostgreSQL extension. The SQL standard leaves the definition of users to the implementation.

See Also
DROP ROLE
DROP USER MAPPING

Name

DROP USER MAPPING — remove a user mapping for a foreign server

Synopsis

DROP USER MAPPING [ IF EXISTS ] FOR { user_name | USER | CURRENT_USER | PUBLIC } SERVER server_name

Description

DROP USER MAPPING removes an existing user mapping from foreign server.

The owner of a foreign server can drop user mappings for that server for any user. Also, a user can drop a user mapping for their own user name if USAGE privilege on the server has been granted to the user.

Parameters

IF EXISTS

Do not throw an error if the user mapping does not exist. A notice is issued in this case.

user_name

User name of the mapping. CURRENT_USER and USER match the name of the current user. PUBLIC is used to match all present and future user names in the system.

server_name

Server name of the user mapping.

Examples

Drop a user mapping bob, server foo if it exists:

DROP USER MAPPING IF EXISTS FOR bob SERVER foo;

Compatibility

DROP USER MAPPING conforms to ISO/IEC 9075-9 (SQL/MED). The IF EXISTS clause is a PostgreSQL extension.
See Also

CREATE USER MAPPING, ALTER USER MAPPING
DROP VIEW

Name
DROP VIEW — remove a view

Synopsis
DROP VIEW [ IF EXISTS ] name [, ...] [ CASCADE | RESTRICT ]

Description
DROP VIEW drops an existing view. To execute this command you must be the owner of the view.

Parameters
IF EXISTS
Do not throw an error if the view does not exist. A notice is issued in this case.
name
The name (optionally schema-qualified) of the view to remove.
CASCADE
Automatically drop objects that depend on the view (such as other views), and in turn all objects that depend on those objects (see Section 5.13).
RESTRICT
Refuse to drop the view if any objects depend on it. This is the default.

Examples
This command will remove the view called kinds:

DROP VIEW kinds;

Compatibility
This command conforms to the SQL standard, except that the standard only allows one view to be dropped per command, and apart from the IF EXISTS option, which is a PostgreSQL extension.

See Also
ALTER VIEW, CREATE VIEW
**Name**

END — commit the current transaction

**Synopsis**

END [ WORK | TRANSACTION ]

**Description**

END commits the current transaction. All changes made by the transaction become visible to others and are guaranteed to be durable if a crash occurs. This command is a PostgreSQL extension that is equivalent to COMMIT.

**Parameters**

WORK
TRANSACTION

Optional key words. They have no effect.

**Notes**

Use ROLLBACK to abort a transaction.

Issuing END when not inside a transaction does no harm, but it will provoke a warning message.

**Examples**

To commit the current transaction and make all changes permanent:

    END;

**Compatibility**

END is a PostgreSQL extension that provides functionality equivalent to COMMIT, which is specified in the SQL standard.

**See Also**

BEGIN, COMMIT, ROLLBACK
EXECUTE

Name
EXECUTE — execute a prepared statement

Synopsis
EXECUTE name [ ( parameter [, ...] ) ]

Description
EXECUTE is used to execute a previously prepared statement. Since prepared statements only exist for the duration of a session, the prepared statement must have been created by a PREPARE statement executed earlier in the current session.

If the PREPARE statement that created the statement specified some parameters, a compatible set of parameters must be passed to the EXECUTE statement, or else an error is raised. Note that (unlike functions) prepared statements are not overloaded based on the type or number of their parameters; the name of a prepared statement must be unique within a database session.

For more information on the creation and usage of prepared statements, see PREPARE.

Parameters

name
The name of the prepared statement to execute.

parameter
The actual value of a parameter to the prepared statement. This must be an expression yielding a value that is compatible with the data type of this parameter, as was determined when the prepared statement was created.

Outputs
The command tag returned by EXECUTE is that of the prepared statement, and not EXECUTE.

Examples
Examples are given in the Examples section of the PREPARE documentation.

Compatibility
The SQL standard includes an EXECUTE statement, but it is only for use in embedded SQL. This version of the EXECUTE statement also uses a somewhat different syntax.
See Also

DEALLOCATE, PREPARE
**EXPLAIN**

**Name**

EXPLAIN — show the execution plan of a statement

**Synopsis**

EXPLAIN [ ( option [, ...] ) ] statement  
EXPLAIN [ ANALYZE ] [ VERBOSE ] statement

where option can be one of:

- ANALYZE [ boolean ]
- VERBOSE [ boolean ]
- COSTS [ boolean ]
- BUFFERS [ boolean ]
- TIMING [ boolean ]
- FORMAT { TEXT | XML | JSON | YAML }

**Description**

This command displays the execution plan that the PostgreSQL planner generates for the supplied statement. The execution plan shows how the table(s) referenced by the statement will be scanned — by plain sequential scan, index scan, etc. — and if multiple tables are referenced, what join algorithms will be used to bring together the required rows from each input table.

The most critical part of the display is the estimated statement execution cost, which is the planner’s guess at how long it will take to run the statement (measured in cost units that are arbitrary, but conventionally mean disk page fetches). Actually two numbers are shown: the start-up cost before the first row can be returned, and the total cost to return all the rows. For most queries the total cost is what matters, but in contexts such as a subquery in EXISTS, the planner will choose the smallest start-up cost instead of the smallest total cost (since the executor will stop after getting one row, anyway). Also, if you limit the number of rows to return with a LIMIT clause, the planner makes an appropriate interpolation between the endpoint costs to estimate which plan is really the cheapest.

The ANALYZE option causes the statement to be actually executed, not only planned. Then actual run time statistics are added to the display, including the total elapsed time expended within each plan node (in milliseconds) and the total number of rows it actually returned. This is useful for seeing whether the planner’s estimates are close to reality.

**Important:** Keep in mind that the statement is actually executed when the ANALYZE option is used. Although EXPLAIN will discard any output that a SELECT would return, other side effects of the statement will happen as usual. If you wish to use EXPLAIN ANALYZE on an INSERT, UPDATE, DELETE, CREATE TABLE AS, or EXECUTE statement without letting the command affect your data, use this approach:

BEGIN;  
EXPLAIN ANALYZE ...;  
ROLLBACK;
Only the `ANALYZE` and `VERBOSE` options can be specified, and only in that order, without surrounding the option list in parentheses. Prior to PostgreSQL 9.0, the unparenthesized syntax was the only one supported. It is expected that all new options will be supported only in the parenthesized syntax.

## Parameters

**ANALYZE**

Carry out the command and show actual run times and other statistics. This parameter defaults to `FALSE`.

**VERBOSE**

Display additional information regarding the plan. Specifically, include the output column list for each node in the plan tree, schema-qualify table and function names, always label variables in expressions with their range table alias, and always print the name of each trigger for which statistics are displayed. This parameter defaults to `FALSE`.

**COSTS**

Include information on the estimated startup and total cost of each plan node, as well as the estimated number of rows and the estimated width of each row. This parameter defaults to `TRUE`.

**BUFFERS**

Include information on buffer usage. Specifically, include the number of shared blocks hit, read, dirtied, and written, the number of local blocks hit, read, dirtied, and written, and the number of temp blocks read and written. A *hit* means that a read was avoided because the block was found already in cache when needed. Shared blocks contain data from regular tables and indexes; local blocks contain data from temporary tables and indexes; while temp blocks contain short-term working data used in sorts, hashes, Materialize plan nodes, and similar cases. The number of blocks *dirtied* indicates the number of previously unmodified blocks that were changed by this query; while the number of blocks *written* indicates the number of previously-dirtied blocks evicted from cache by this backend during query processing. The number of blocks shown for an upper-level node includes those used by all its child nodes. In text format, only non-zero values are printed. This parameter may only be used when `ANALYZE` is also enabled. It defaults to `FALSE`.

**TIMING**

Include actual startup time and time spent in each node in the output. The overhead of repeatedly reading the system clock can slow down the query significantly on some systems, so it may be useful to set this parameter to `FALSE` when only actual row counts, and not exact times, are needed. Run time of the entire statement is always measured, even when node-level timing is turned off with this option. This parameter may only be used when `ANALYZE` is also enabled. It defaults to `TRUE`.

**FORMAT**

Specify the output format, which can be `TEXT`, `XML`, `JSON`, or `YAML`. Non-text output contains the same information as the text output format, but is easier for programs to parse. This parameter defaults to `TEXT`. 

1646
**EXPLAIN**

**boolean**

Specifies whether the selected option should be turned on or off. You can write `TRUE`, `ON`, or `1` to enable the option, and `FALSE`, `OFF`, or `0` to disable it. The boolean value can also be omitted, in which case `TRUE` is assumed.

**statement**

Any `SELECT`, `INSERT`, `UPDATE`, `DELETE`, `VALUES`, `EXECUTE`, `DECLARE`, `CREATE TABLE AS`, or `CREATE MATERIALIZED VIEW AS` statement, whose execution plan you wish to see.

**Outputs**

The command’s result is a textual description of the plan selected for the statement, optionally annotated with execution statistics. Section 14.1 describes the information provided.

**Notes**

In order to allow the PostgreSQL query planner to make reasonably informed decisions when optimizing queries, the `pg_statistic` data should be up-to-date for all tables used in the query. Normally the autovacuum daemon will take care of that automatically. But if a table has recently had substantial changes in its contents, you might need to do a manual `ANALYZE` rather than wait for autovacuum to catch up with the changes.

In order to measure the run-time cost of each node in the execution plan, the current implementation of `EXPLAIN ANALYZE` adds profiling overhead to query execution. As a result, running `EXPLAIN ANALYZE` on a query can sometimes take significantly longer than executing the query normally. The amount of overhead depends on the nature of the query, as well as the platform being used. The worst case occurs for plan nodes that in themselves require very little time per execution, and on machines that have relatively slow operating system calls for obtaining the time of day.

**Examples**

To show the plan for a simple query on a table with a single `integer` column and 10000 rows:

```
EXPLAIN SELECT * FROM foo;
```

```
QUERY PLAN
---------------------------------------------------------
Seq Scan on foo (cost=0.00..155.00 rows=10000 width=4)
(1 row)
```

Here is the same query, with JSON output formatting:

```
EXPLAIN (FORMAT JSON) SELECT * FROM foo;
```

```
[{
  "Plan": {
    "Node Type": "Seq Scan",
    "Relation Name": "foo",
    "Startup Cost": 0,
    "Estimated Cost": 155.00
  }
}]
```

1647
If there is an index and we use a query with an indexable WHERE condition, EXPLAIN might show a different plan:

```
EXPLAIN SELECT * FROM foo WHERE i = 4;
```

```
QUERY PLAN
-----------------------------------------------------------------------------
Index Scan using fi on foo (cost=0.00..5.98 rows=1 width=4)
  Index Cond: (i = 4)
(2 rows)
```

Here is the same query, but in YAML format:

```
EXPLAIN (FORMAT YAML) SELECT * FROM foo WHERE i='4';
```

```
- Plan:
  - Node Type: "Index Scan"
  - Scan Direction: "Forward"
  - Index Name: "fi"
  - Relation Name: "foo"
  - Alias: "foo"
  - Startup Cost: 0.00
  - Total Cost: 5.98
  - Plan Rows: 1
  - Plan Width: 4
  - Index Cond: "(i = 4)"

(1 row)
```

XML format is left as an exercise for the reader.

Here is the same plan with cost estimates suppressed:

```
EXPLAIN (COSTS FALSE) SELECT * FROM foo WHERE i = 4;
```

```
QUERY PLAN
-----------------------------------------------------------------------------
Index Scan using fi on foo
  Index Cond: (i = 4)
(2 rows)
```

Here is an example of a query plan for a query using an aggregate function:

```
EXPLAIN SELECT sum(i) FROM foo WHERE i < 10;
```
Here is an example of using `EXPLAIN EXECUTE` to display the execution plan for a prepared query:

```sql
PREPARE query(int, int) AS SELECT sum(bar) FROM test
  WHERE id > $1 AND id < $2
  GROUP BY foo;

EXPLAIN ANALYZE EXECUTE query(100, 200);
```

Of course, the specific numbers shown here depend on the actual contents of the tables involved. Also note that the numbers, and even the selected query strategy, might vary between PostgreSQL releases due to planner improvements. In addition, the `ANALYZE` command uses random sampling to estimate data statistics; therefore, it is possible for cost estimates to change after a fresh run of `ANALYZE`, even if the actual distribution of data in the table has not changed.

**Compatibility**

There is no `EXPLAIN` statement defined in the SQL standard.

**See Also**

`ANALYZE`
FETCH

Name

FETCH — retrieve rows from a query using a cursor

Synopsis

FETCH [ direction [ FROM | IN ] ] cursor_name

where direction can be empty or one of:

NEXT
PRIOR
FIRST
LAST
ABSOLUTE count
RELATIVE count
ALL
FORWARD
FORWARD count
FORWARD ALL
BACKWARD
BACKWARD count
BACKWARD ALL

Description

FETCH retrieves rows using a previously-created cursor.

A cursor has an associated position, which is used by FETCH. The cursor position can be before the first row of the query result, on any particular row of the result, or after the last row of the result. When created, a cursor is positioned before the first row. After fetching some rows, the cursor is positioned on the row most recently retrieved. If FETCH runs off the end of the available rows then the cursor is left positioned after the last row, or before the first row if fetching backward. FETCH ALL or FETCH BACKWARD ALL will always leave the cursor positioned after the last row or before the first row.

The forms NEXT, PRIOR, FIRST, LAST, ABSOLUTE, RELATIVE fetch a single row after moving the cursor appropriately. If there is no such row, an empty result is returned, and the cursor is left positioned before the first row or after the last row as appropriate.

The forms using FORWARD and BACKWARD retrieve the indicated number of rows moving in the forward or backward direction, leaving the cursor positioned on the last-returned row (or after/before all rows, if the count exceeds the number of rows available).

RELATIVE 0, FORWARD 0, and BACKWARD 0 all request fetching the current row without moving the cursor, that is, re-fetching the most recently fetched row. This will succeed unless the cursor is positioned before the first row or after the last row; in which case, no row is returned.

Note: This page describes usage of cursors at the SQL command level. If you are trying to use cursors inside a PL/pgSQL function, the rules are different — see Section 41.7.3.
Parameters

direction

direction defines the fetch direction and number of rows to fetch. It can be one of the following:

NEXT
   Fetch the next row. This is the default if direction is omitted.

PRIOR
   Fetch the prior row.

FIRST
   Fetch the first row of the query (same as ABSOLUTE 1).

LAST
   Fetch the last row of the query (same as ABSOLUTE -1).

ABSOLUTE count
   Fetch the count’th row of the query, or the abs(count)’th row from the end if count is negative. Position before first row or after last row if count is out of range; in particular, ABSOLUTE 0 positions before the first row.

RELATIVE count
   Fetch the count’th succeeding row, or the abs(count)’th prior row if count is negative. RELATIVE 0 re-fetches the current row, if any.

count
   Fetch the next count rows (same as FORWARD count).

ALL
   Fetch all remaining rows (same as FORWARD ALL).

FORWARD
   Fetch the next row (same as NEXT).

FORWARD count
   Fetch the next count rows. FORWARD 0 re-fetches the current row.

FORWARD ALL
   Fetch all remaining rows.

BACKWARD
   Fetch the prior row (same as PRIOR).

BACKWARD count
   Fetch the prior count rows (scanning backwards). BACKWARD 0 re-fetches the current row.

BACKWARD ALL
   Fetch all prior rows (scanning backwards).
**FETCH**

count

count is a possibly-signed integer constant, determining the location or number of rows to fetch. For FORWARD and BACKWARD cases, specifying a negative count is equivalent to changing the sense of FORWARD and BACKWARD.

cursor_name

An open cursor’s name.

**Outputs**

On successful completion, a FETCH command returns a command tag of the form

FETCH count

The count is the number of rows fetched (possibly zero). Note that in psql, the command tag will not actually be displayed, since psql displays the fetched rows instead.

**Notes**

The cursor should be declared with the SCROLL option if one intends to use any variants of FETCH other than FETCH NEXT or FETCH FORWARD with a positive count. For simple queries PostgreSQL will allow backwards fetch from cursors not declared with SCROLL, but this behavior is best not relied on. If the cursor is declared with NO SCROLL, no backward fetches are allowed.

ABSOLUTE fetches are not any faster than navigating to the desired row with a relative move: the underlying implementation must traverse all the intermediate rows anyway. Negative absolute fetches are even worse: the query must be read to the end to find the last row, and then traversed backward from there. However, rewinding to the start of the query (as with FETCH ABSOLUTE 0) is fast.

DECLARE is used to define a cursor. Use MOVE to change cursor position without retrieving data.

**Examples**

The following example traverses a table using a cursor:

```
BEGIN WORK;

-- Set up a cursor:
DECLARE liahona SCROLL CURSOR FOR SELECT * FROM films;

-- Fetch the first 5 rows in the cursor liahona:
FETCH FORWARD 5 FROM liahona;

<table>
<thead>
<tr>
<th>code</th>
<th>title</th>
<th>did</th>
<th>date_prod</th>
<th>kind</th>
<th>len</th>
</tr>
</thead>
<tbody>
<tr>
<td>BL101</td>
<td>The Third Man</td>
<td>101</td>
<td>1949-12-23</td>
<td>Drama</td>
<td>01:44</td>
</tr>
<tr>
<td>BL102</td>
<td>The African Queen</td>
<td>101</td>
<td>1951-08-11</td>
<td>Romantic</td>
<td>01:43</td>
</tr>
<tr>
<td>JL201</td>
<td>Une Femme est une Femme</td>
<td>102</td>
<td>1961-03-12</td>
<td>Romantic</td>
<td>01:25</td>
</tr>
<tr>
<td>P_301</td>
<td>Vertigo</td>
<td>103</td>
<td>1958-11-14</td>
<td>Action</td>
<td>02:08</td>
</tr>
<tr>
<td>P_302</td>
<td>Becket</td>
<td>103</td>
<td>1964-02-03</td>
<td>Drama</td>
<td>02:28</td>
</tr>
</tbody>
</table>

-- Fetch the previous row:
```
FETCH PRIOR FROM liahona;

<table>
<thead>
<tr>
<th>code</th>
<th>title</th>
<th>did</th>
<th>date_prod</th>
<th>kind</th>
<th>len</th>
</tr>
</thead>
<tbody>
<tr>
<td>P_301</td>
<td>Vertigo</td>
<td>103</td>
<td>1958-11-14</td>
<td>Action</td>
<td>02:08</td>
</tr>
</tbody>
</table>

-- Close the cursor and end the transaction:
CLOSE liahona;
COMMIT WORK;

Compatibility

The SQL standard defines FETCH for use in embedded SQL only. The variant of FETCH described here returns the data as if it were a SELECT result rather than placing it in host variables. Other than this point, FETCH is fully upward-compatible with the SQL standard.

The FETCH forms involving FORWARD and BACKWARD, as well as the forms FETCH count and FETCH ALL, in which FORWARD is implicit, are PostgreSQL extensions.

The SQL standard allows only FROM preceding the cursor name; the option to use IN, or to leave them out altogether, is an extension.

See Also

CLOSE, DECLARE, MOVE
GRANT

Name
GRANT — define access privileges

Synopsis

GRANT { { SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES | TRIGGER } [ , ... ] | ALL [ PRIVILEGES ] } ON { [ TABLE ] table_name [ , ... ] | ALL TABLES IN SCHEMA schema_name [ , ... ] } TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { { SELECT | INSERT | UPDATE | REFERENCES } ( column_name [ , ... ] ) [ , ... ] | ALL [ PRIVILEGES ] ( column_name [ , ... ] ) } ON [ TABLE ] table_name [ , ... ] TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { { USAGE | SELECT | UPDATE } [ , ... ] | ALL [ PRIVILEGES ] } ON { SEQUENCE sequence_name [ , ... ] | ALL SEQUENCES IN SCHEMA schema_name [ , ... ] } TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { { CREATE | CONNECT | TEMPORARY | TEMP } [ , ... ] | ALL [ PRIVILEGES ] } ON DATABASE database_name [ , ... ] TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] } ON DOMAIN domain_name [ , ... ] TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] } ON FOREIGN DATA WRAPPER fdw_name [ , ... ] TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] } ON FOREIGN SERVER server_name [ , ... ] TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { EXECUTE | ALL [ PRIVILEGES ] } ON { FUNCTION function_name ( [ [ argmode ] [ arg_name ] arg_type [ , ... ] ] ) [ , ... ] | ALL FUNCTIONS IN SCHEMA schema_name [ , ... ] } TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] } ON LANGUAGE lang_name [ , ... ] TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { { SELECT | UPDATE } [ , ... ] | ALL [ PRIVILEGES ] } ON LARGE OBJECTloid [ , ... ] TO role_specification [ , ... ] [ WITH GRANT OPTION ]

GRANT { { CREATE | USAGE } [ , ... ] | ALL [ PRIVILEGES ] }
GRANT

ON SCHEMA schema_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { CREATE | ALL [ PRIVILEGES ] }
ON TABLESPACE tablespace_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

GRANT { USAGE | ALL [ PRIVILEGES ] }
ON TYPE type_name [, ...]
TO role_specification [, ...] [ WITH GRANT OPTION ]

where role_specification can be:

[ GROUP ] role_name
| PUBLIC
| CURRENT_USER
| SESSION_USER

GRANT role_name [, ...] TO role_name [, ...] [ WITH ADMIN OPTION ]

Description

The GRANT command has two basic variants: one that grants privileges on a database object (table, column, view, foreign table, sequence, database, foreign-data wrapper, foreign server, function, procedural language, schema, or tablespace), and one that grants membership in a role. These variants are similar in many ways, but they are different enough to be described separately.

GRANT on Database Objects

This variant of the GRANT command gives specific privileges on a database object to one or more roles. These privileges are added to those already granted, if any.

There is also an option to grant privileges on all objects of the same type within one or more schemas. This functionality is currently supported only for tables, sequences, and functions (but note that ALL TABLES is considered to include views and foreign tables).

The key word PUBLIC indicates that the privileges are to be granted to all roles, including those that might be created later. PUBLIC can be thought of as an implicitly defined group that always includes all roles. Any particular role will have the sum of privileges granted directly to it, privileges granted to any role it is presently a member of, and privileges granted to PUBLIC.

If WITH GRANT OPTION is specified, the recipient of the privilege can in turn grant it to others. Without a grant option, the recipient cannot do that. Grant options cannot be granted to PUBLIC.

There is no need to grant privileges to the owner of an object (usually the user that created it), as the owner has all privileges by default. (The owner could, however, choose to revoke some of their own privileges for safety.)

The right to drop an object, or to alter its definition in any way, is not treated as a grantable privilege; it is inherent in the owner, and cannot be granted or revoked. (However, a similar effect can be obtained by granting or revoking membership in the role that owns the object; see below.) The owner implicitly has all grant options for the object, too.

PostgreSQL grants default privileges on some types of objects to PUBLIC. No privileges are granted to PUBLIC by default on tables, table columns, sequences, foreign data wrappers, foreign servers, large objects, schemas, or tablespaces. For other types of objects, the default privileges granted to PUBLIC
GRANT

are as follows: CONNECT and TEMPORARY (create temporary tables) privileges for databases; EXECUTE privilege for functions; and USAGE privilege for languages and data types (including domains). The object owner can, of course, REVOKE both default and expressly granted privileges. (For maximum security, issue the REVOKE in the same transaction that creates the object; then there is no window in which another user can use the object.) Also, these initial default privilege settings can be changed using the ALTER DEFAULT PRIVILEGES command.

The possible privileges are:

SELECT

Allows SELECT from any column, or the specific columns listed, of the specified table, view, or sequence. Also allows the use of COPY TO. This privilege is also needed to reference existing column values in UPDATE or DELETE. For sequences, this privilege also allows the use of the currval function. For large objects, this privilege allows the object to be read.

INSERT

Allows INSERT of a new row into the specified table. If specific columns are listed, only those columns may be assigned to in the INSERT command (other columns will therefore receive default values). Also allows COPY FROM.

UPDATE

Allows UPDATE of any column, or the specific columns listed, of the specified table. (In practice, any nontrivial UPDATE command will require SELECT privilege as well, since it must reference table columns to determine which rows to update, and/or to compute new values for columns.) SELECT ... FOR UPDATE and SELECT ... FOR SHARE also require this privilege on at least one column, in addition to the SELECT privilege. For sequences, this privilege allows the use of the nextval and setval functions. For large objects, this privilege allows writing or truncating the object.

DELETE

Allows DELETE of a row from the specified table. (In practice, any nontrivial DELETE command will require SELECT privilege as well, since it must reference table columns to determine which rows to delete.)

TRUNCATE

Allows TRUNCATE on the specified table.

REFERENCES

To create a foreign key constraint, it is necessary to have this privilege on both the referencing and referenced columns. The privilege may be granted for all columns of a table, or just specific columns.

TRIGGER

Allows the creation of a trigger on the specified table. (See the CREATE TRIGGER statement.)

CREATE

For databases, allows new schemas to be created within the database.

For schemas, allows new objects to be created within the schema. To rename an existing object, you must own the object and have this privilege for the containing schema.

For tablespaces, allows tables, indexes, and temporary files to be created within the tablespace, and allows databases to be created that have the tablespace as their default tablespace. (Note that revoking this privilege will not alter the placement of existing objects.)
CONNECT

Allows the user to connect to the specified database. This privilege is checked at connection startup (in addition to checking any restrictions imposed by pg_hba.conf).

TEMPORARY

TEMP

Allows temporary tables to be created while using the specified database.

EXECUTE

Allows the use of the specified function and the use of any operators that are implemented on top of the function. This is the only type of privilege that is applicable to functions. (This syntax works for aggregate functions, as well.)

USAGE

For procedural languages, allows the use of the specified language for the creation of functions in that language. This is the only type of privilege that is applicable to procedural languages.

For schemas, allows access to objects contained in the specified schema (assuming that the objects’ own privilege requirements are also met). Essentially this allows the grantee to “look up” objects within the schema. Without this permission, it is still possible to see the object names, e.g. by querying the system tables. Also, after revoking this permission, existing backends might have statements that have previously performed this lookup, so this is not a completely secure way to prevent object access.

For sequences, this privilege allows the use of the currval and nextval functions.

For types and domains, this privilege allow the use of the type or domain in the creation of tables, functions, and other schema objects. (Note that it does not control general “usage” of the type, such as values of the type appearing in queries. It only prevents objects from being created that depend on the type. The main purpose of the privilege is controlling which users create dependencies on a type, which could prevent the owner from changing the type later.)

For foreign-data wrappers, this privilege enables the grantee to create new servers using that foreign-data wrapper.

For servers, this privilege enables the grantee to create foreign tables using the server, and also to create, alter, or drop their own user’s user mappings associated with that server.

ALL PRIVILEGES

Grant all of the available privileges at once. The PRIVILEGES key word is optional in PostgreSQL, though it is required by strict SQL.

The privileges required by other commands are listed on the reference page of the respective command.

GRANT on Roles

This variant of the GRANT command grants membership in a role to one or more other roles. Membership in a role is significant because it conveys the privileges granted to a role to each of its members.

If WITH ADMIN OPTION is specified, the member can in turn grant membership in the role to others, and revoke membership in the role as well. Without the admin option, ordinary users cannot do that. A role is not considered to hold WITH ADMIN OPTION on itself, but it may grant or revoke membership in itself from a database session where the session user matches the role. Database superusers can grant or revoke membership in any role to anyone. Roles having CREATEROLE privilege can grant or revoke membership in any role that is not a superuser.
Unlike the case with privileges, membership in a role cannot be granted to `PUBLIC`. Note also that this form of the command does not allow the noise word `GROUP`.

### Notes

The `REVOKE` command is used to revoke access privileges.

Since PostgreSQL 8.1, the concepts of users and groups have been unified into a single kind of entity called a role. It is therefore no longer necessary to use the keyword `GROUP` to identify whether a grantee is a user or a group. `GROUP` is still allowed in the command, but it is a noise word.

A user may perform `SELECT`, `INSERT`, etc. on a column if they hold that privilege for either the specific column or its whole table. Granting the privilege at the table level and then revoking it for one column will not do what one might wish: the table-level grant is unaffected by a column-level operation.

When a non-owner of an object attempts to `GRANT` privileges on the object, the command will fail outright if the user has no privileges whatsoever on the object. As long as some privilege is available, the command will proceed, but it will grant only those privileges for which the user has grant options. The `GRANT ALL PRIVILEGES` forms will issue a warning message if no grant options are held, while the other forms will issue a warning if grant options for any of the privileges specifically named in the command are not held. (In principle these statements apply to the object owner as well, but since the owner is always treated as holding all grant options, the cases can never occur.)

It should be noted that database superusers can access all objects regardless of object privilege settings. This is comparable to the rights of `root` in a Unix system. As with `root`, it’s unwise to operate as a superuser except when absolutely necessary.

If a superuser chooses to issue a `GRANT` or `REVOKE` command, the command is performed as though it were issued by the owner of the affected object. In particular, privileges granted via such a command will appear to have been granted by the object owner. (For role membership, the membership appears to have been granted by the containing role itself.)

`GRANT` and `REVOKE` can also be done by a role that is not the owner of the affected object, but is a member of the role that owns the object, or is a member of a role that holds privileges `WITH GRANT OPTION` on the object. In this case the privileges will be recorded as having been granted by the role that actually owns the object or holds the privileges `WITH GRANT OPTION`. For example, if table `t1` is owned by role `g1`, of which role `u1` is a member, then `u1` can grant privileges on `t1` to `u2`, but those privileges will appear to have been granted directly by `g1`. Any other member of role `g1` could revoke them later.

If the role executing `GRANT` holds the required privileges indirectly via more than one role membership path, it is unspecified which containing role will be recorded as having done the grant. In such cases it is best practice to use `SET ROLE` to become the specific role you want to do the `GRANT` as.

Granting permission on a table does not automatically extend permissions to any sequences used by the table, including sequences tied to `SERIAL` columns. Permissions on sequences must be set separately.

Use `psql`’s ```dp``` command to obtain information about existing privileges for tables and columns. For example:

```bash
=> \dp mytable

<table>
<thead>
<tr>
<th>Schema</th>
<th>Name</th>
<th>Type</th>
<th>Access privileges</th>
<th>Column access privileges</th>
</tr>
</thead>
</table>
```
The entries shown by \dp are interpreted thus:

rolename=xxxx -- privileges granted to a role
=xxxx -- privileges granted to PUBLIC

\[\begin{array}{ll}
\text{r} & \text{SELECT } ("\text{read}\") \\
\text{w} & \text{UPDATE } ("\text{write}\") \\
\text{a} & \text{INSERT } ("\text{append}\") \\
\text{d} & \text{DELETE} \\
\text{D} & \text{TRUNCATE} \\
\text{x} & \text{REFERENCES} \\
\text{t} & \text{TRIGGER} \\
\text{X} & \text{EXECUTE} \\
\text{U} & \text{USAGE} \\
\text{C} & \text{CREATE} \\
\text{c} & \text{CONNECT} \\
\text{T} & \text{TEMPORARY} \\
arwdDxt & \text{ALL PRIVILEGES (for tables, varies for other objects)} \\
\ast & \text{grant option for preceding privilege} \\
\text{/yyyy} & \text{role that granted this privilege}
\end{array}\]

The above example display would be seen by user \texttt{miriam} after creating table \texttt{mytable} and doing:

GRANT SELECT ON mytable TO PUBLIC;
GRANT SELECT, UPDATE, INSERT ON mytable TO admin;
GRANT SELECT (col1), UPDATE (col1) ON mytable TO miriam_rw;

For non-table objects there are other \d commands that can display their privileges.

If the “Access privileges” column is empty for a given object, it means the object has default privileges (that is, its privileges column is null). Default privileges always include all privileges for the owner, and can include some privileges for PUBLIC depending on the object type, as explained above. The first GRANT or REVOKE on an object will instantiate the default privileges (producing, for example, \{miriam\}=arwdDxt/miriam) and then modify them per the specified request. Similarly, entries are shown in “Column access privileges” only for columns with nondefault privileges. (Note: for this purpose, “default privileges” always means the built-in default privileges for the object’s type. An object whose privileges have been affected by an ALTER DEFAULT PRIVILEGES command will always be shown with an explicit privilege entry that includes the effects of the ALTER.)

Notice that the owner’s implicit grant options are not marked in the access privileges display. A * will appear only when grant options have been explicitly granted to someone.

Examples

Grant insert privilege to all users on table films:

GRANT INSERT ON films TO PUBLIC;
Grant all available privileges to user manuel on view kinds:

GRANT ALL PRIVILEGES ON kinds TO manuel;

Note that while the above will indeed grant all privileges if executed by a superuser or the owner of kinds, when executed by someone else it will only grant those permissions for which the someone else has grant options.

Grant membership in role admins to user joe:

GRANT admins TO joe;

Compatibility

According to the SQL standard, the PRIVILEGES key word in ALL PRIVILEGES is required. The SQL standard does not support setting the privileges on more than one object per command.

PostgreSQL allows an object owner to revoke their own ordinary privileges: for example, a table owner can make the table read-only to themselves by revoking their own INSERT, UPDATE, DELETE, and TRUNCATE privileges. This is not possible according to the SQL standard. The reason is that PostgreSQL treats the owner’s privileges as having been granted by the owner to themselves; therefore they can revoke them too. In the SQL standard, the owner’s privileges are granted by an assumed entity “_SYSTEM”. Not being “_SYSTEM”, the owner cannot revoke these rights.

According to the SQL standard, grant options can be granted to PUBLIC; PostgreSQL only supports granting grant options to roles.

The SQL standard provides for a USAGE privilege on other kinds of objects: character sets, collations, translations.

In the SQL standard, sequences only have a USAGE privilege, which controls the use of the NEXT VALUE FOR expression, which is equivalent to the function nextval in PostgreSQL. The sequence privileges SELECT and UPDATE are PostgreSQL extensions. The application of the sequence USAGE privilege to the currval function is also a PostgreSQL extension (as is the function itself).

Privileges on databases, tablespaces, schemas, and languages are PostgreSQL extensions.

See Also

REVOKE, ALTER DEFAULT PRIVILEGES
IMPORT FOREIGN SCHEMA

Name
IMPORT FOREIGN SCHEMA — import table definitions from a foreign server

Synopsis
IMPORT FOREIGN SCHEMA remote_schema
[ { LIMIT TO | EXCEPT } ( table_name [, ...] ) ]
FROM SERVER server_name
INTO local_schema
[ OPTIONS ( option 'value' [, ... ] ) ]

Description
IMPORT FOREIGN SCHEMA creates foreign tables that represent tables existing on a foreign server. The new foreign tables will be owned by the user issuing the command and are created with the correct column definitions and options to match the remote tables.

By default, all tables and views existing in a particular schema on the foreign server are imported. Optionally, the list of tables can be limited to a specified subset, or specific tables can be excluded. The new foreign tables are all created in the target schema, which must already exist.

To use IMPORT FOREIGN SCHEMA, the user must have USAGE privilege on the foreign server, as well as CREATE privilege on the target schema.

Parameters
remote_schema
The remote schema to import from. The specific meaning of a remote schema depends on the foreign data wrapper in use.

LIMIT TO ( table_name [, ... ] )
Import only foreign tables matching one of the given table names. Other tables existing in the foreign schema will be ignored.

EXCEPT ( table_name [, ... ] )
Exclude specified foreign tables from the import. All tables existing in the foreign schema will be imported except the ones listed here.

server_name
The foreign server to import from.

local_schema
The schema in which the imported foreign tables will be created.
OPTIONS { option 'value' [, ...] }

Options to be used during the import. The allowed option names and values are specific to each foreign data wrapper.

Examples

Import table definitions from a remote schema foreign_films on server film_server, creating the foreign tables in local schema films:

```
IMPORT FOREIGN SCHEMA foreign_films
    FROM SERVER film_server INTO films;
```

As above, but import only the two tables actors and directors (if they exist):

```
IMPORT FOREIGN SCHEMA foreign_films LIMIT TO (actors, directors)
    FROM SERVER film_server INTO films;
```

Compatibility

The IMPORT FOREIGN SCHEMA command conforms to the SQL standard, except that the OPTIONS clause is a PostgreSQL extension.

See Also

CREATE FOREIGN TABLE, CREATE SERVER
INSERT

Name
INSERT — create new rows in a table

Synopsis

[ WITH [ RECURSIVE ] with_query [, ...] ]

INSERT INTO table_name [ AS alias ] [ ( column_name [, ...] ) ]
{ DEFAULT VALUES | VALUES ( { expression | DEFAULT } [, ...] ) [, ...] | query }
[ ON CONFLICT [ conflict_target ] conflict_action ]
[ RETURNING * | output_expression [ [ AS ] output_name ] [, ...] ]

where conflict_target can be one of:

( { index_column_name | ( index_expression ) } [ COLLATE collation ] [ opclass ] [, ...])
ON CONSTRAINT constraint_name

and conflict_action is one of:

DO NOTHING
DO UPDATE SET { column_name = { expression | DEFAULT } |
( column_name [, ...] ) = ( { expression | DEFAULT } [, ...] ) |
( column_name [, ...] ) = ( sub-SELECT )
} [, ...]
[ WHERE condition ]

Description

INSERT inserts new rows into a table. One can insert one or more rows specified by value expressions, or zero or more rows resulting from a query.

The target column names can be listed in any order. If no list of column names is given at all, the default is all the columns of the table in their declared order; or the first \( N \) column names, if there are only \( N \) columns supplied by the VALUES clause or query. The values supplied by the VALUES clause or query are associated with the explicit or implicit column list left-to-right.

Each column not present in the explicit or implicit column list will be filled with a default value, either its declared default value or null if there is none.

If the expression for any column is not of the correct data type, automatic type conversion will be attempted.

ON CONFLICT can be used to specify an alternative action to raising a unique constraint or exclusion constraint violation error. (See ON CONFLICT Clause below.)

The optional RETURNING clause causes INSERT to compute and return value(s) based on each row actually inserted (or updated, if an ON CONFLICT DO UPDATE clause was used). This is primarily useful for obtaining values that were supplied by defaults, such as a serial sequence number. However, any expression using the table’s columns is allowed. The syntax of the RETURNING list is identical to that of the output list of SELECT. Only rows that were successfully inserted or updated will be returned. For example, if a row was locked but not updated because an ON CONFLICT DO UPDATE ... WHERE clause condition was not satisfied, the row will not be returned.
You must have \texttt{INSERT} privilege on a table in order to insert into it. If \texttt{ON CONFLICT DO UPDATE} is present, \texttt{UPDATE} privilege on the table is also required.

If a column list is specified, you only need \texttt{INSERT} privilege on the listed columns. Similarly, when \texttt{ON CONFLICT DO UPDATE} is specified, you only need \texttt{UPDATE} privilege on the column(s) that are listed to be updated. However, there also requires \texttt{SELECT} privilege on any column whose values are read in the \texttt{ON CONFLICT DO UPDATE} expressions or condition.

Use of the \texttt{RETURNING} clause requires \texttt{SELECT} privilege on all columns mentioned in \texttt{RETURNING}. If you use the \texttt{query} clause to insert rows from a query, you of course need to have \texttt{SELECT} privilege on any table or column used in the query.

\section*{Parameters}

\subsection*{Inserting}

This section covers parameters that may be used when only inserting new rows. Parameters \textit{exclusively} used with the \texttt{ON CONFLICT} clause are described separately.

\begin{description}
\item[\texttt{with\_query}]
\begin{description}
\item[The \texttt{WITH} clause allows you to specify one or more subqueries that can be referenced by name in the \texttt{INSERT} query. See Section 7.8 and \texttt{SELECT} for details.]
\end{description}
\item[It is possible for the \texttt{query} (\texttt{SELECT} statement) to also contain a \texttt{WITH} clause. In such a case both sets of \texttt{with\_query} can be referenced within the \texttt{query}, but the second one takes precedence since it is more closely nested.]
\end{description}

\item[\texttt{table\_name}]
\begin{description}
\item[The name (optionally schema-qualified) of an existing table.]
\end{description}

\item[\texttt{alias}]
\begin{description}
\item[A substitute name for \texttt{table\_name}. When an alias is provided, it completely hides the actual name of the table. This is particularly useful when \texttt{ON CONFLICT DO UPDATE} targets a table named \texttt{excluded}, since that’s also the name of the special table representing rows proposed for insertion.]
\end{description}

\item[\texttt{column\_name}]
\begin{description}
\item[The name of a column in the table named by \texttt{table\_name}. The column name can be qualified with a subfield name or array subscript, if needed. (Inserting into only some fields of a composite column leaves the other fields null.) When referencing a column with \texttt{ON CONFLICT DO UPDATE}, do not include the table’s name in the specification of a target column. For example, \texttt{INSERT INTO table\_name ... ON CONFLICT DO UPDATE SET table\_name\.col = 1} is invalid (this follows the general behavior for \texttt{UPDATE}).]
\end{description}

\item[\texttt{DEFAULT VALUES}]
\begin{description}
\item[All columns will be filled with their default values.]
\end{description}

\item[\texttt{expression}]
\begin{description}
\item[An expression or value to assign to the corresponding column.]
\end{description}

\item[\texttt{DEFAULT}]
\begin{description}
\item[The corresponding column will be filled with its default value.]
\end{description}

1664
query

A query (SELECT statement) that supplies the rows to be inserted. Refer to the SELECT statement for a description of the syntax.

output_expression

An expression to be computed and returned by the INSERT command after each row is inserted or updated. The expression can use any column names of the table named by table_name. Write * to return all columns of the inserted or updated row(s).

output_name

A name to use for a returned column.

**ON CONFLICT Clause**

The optional ON CONFLICT clause specifies an alternative action to raising a unique violation or exclusion constraint violation error. For each individual row proposed for insertion, either the insertion proceeds, or, if an arbiter constraint or index specified by conflict_target is violated, the alternative conflict_action is taken. ON CONFLICT DO NOTHING simply avoids inserting a row as its alternative action. ON CONFLICT DO UPDATE updates the existing row that conflicts with the row proposed for insertion as its alternative action.

conflict_target can perform unique index inference. When performing inference, it consists of one or more index_column_name columns and/or index_expression expressions, and an optional index_predicate. All table_name unique indexes that, without regard to order, contain exactly the conflict_target-specified columns/expressions are inferred (chosen) as arbiter indexes. If an index_predicate is specified, it must, as a further requirement for inference, satisfy arbiter indexes. Note that this means a non-partial unique index (a unique index without a predicate) will be inferred (and thus used by ON CONFLICT) if such an index satisfying every other criteria is available. If an attempt at inference is unsuccessful, an error is raised.

ON CONFLICT DO UPDATE guarantees an atomic INSERT or UPDATE outcome; provided there is no independent error, one of those two outcomes is guaranteed, even under high concurrency. This is also known as UPSERT — “UPDATE or INSERT”.

conflict_target

Specifies which conflicts ON CONFLICT takes the alternative action on by choosing arbiter indexes. Either performs unique index inference, or names a constraint explicitly. For ON CONFLICT DO NOTHING, it is optional to specify a conflict_target; when omitted, conflicts with all usable constraints (and unique indexes) are handled. For ON CONFLICT DO UPDATE, a conflict_target must be provided.

conflict_action

conflict_action specifies an alternative ON CONFLICT action. It can be either DO NOTHING, or a DO UPDATE clause specifying the exact details of the UPDATE action to be performed in case of a conflict. The SET and WHERE clauses in ON CONFLICT DO UPDATE have access to the existing row using the table’s name (or an alias), and to rows proposed for insertion using the special excluded table. SELECT privilege is required on any column in the target table where corresponding excluded columns are read.

Note that the effects of all per-row BEFORE INSERT triggers are reflected in excluded values, since those effects may have contributed to the row being excluded from insertion.
index_column_name

The name of a table_name column. Used to infer arbiter indexes. Follows CREATE INDEX format. SELECT privilege on index_column_name is required.

index_expression

Similar to index_column_name, but used to infer expressions on table_name columns appearing within index definitions (not simple columns). Follows CREATE INDEX format. SELECT privilege on any column appearing within index_expression is required.

collation

When specified, mandates that corresponding index_column_name or index_expression use a particular collation in order to be matched during inference. Typically this is omitted, as collations usually do not affect whether or not a constraint violation occurs. Follows CREATE INDEX format.

opclass

When specified, mandates that corresponding index_column_name or index_expression use particular operator class in order to be matched during inference. Typically this is omitted, as the equality semantics are often equivalent across a type’s operator classes anyway, or because it’s sufficient to trust that the defined unique indexes have the pertinent definition of equality. Follows CREATE INDEX format.

index_predicate

Used to allow inference of partial unique indexes. Any indexes that satisfy the predicate (which need not actually be partial indexes) can be inferred. Follows CREATE INDEX format. SELECT privilege on any column appearing within index_predicate is required.

constraint_name

Explicitly specifies an arbiter constraint by name, rather than inferring a constraint or index.

condition

An expression that returns a value of type boolean. Only rows for which this expression returns true will be updated, although all rows will be locked when the ON CONFLICT DO UPDATE action is taken. Note that condition is evaluated last, after a conflict has been identified as a candidate to update.

Note that exclusion constraints are not supported as arbiters with ON CONFLICT DO UPDATE. In all cases, only NOT DEFERRABLE constraints and unique indexes are supported as arbiters.

**Tip:** It is often preferable to use unique index inference rather than naming a constraint directly using ON CONFLICT ON CONSTRAINT constraint_name. Inference will continue to work correctly when the underlying index is replaced by another more or less equivalent index in an overlapping way, for example when using CREATE UNIQUE INDEX ... CONCURRENTLY before dropping the index being replaced.
Outputs

On successful completion, an `INSERT` command returns a command tag of the form

```
INSERT oid count
```

The `count` is the number of rows inserted or updated. If `count` is exactly one, and the target table has OIDs, then `oid` is the OID assigned to the inserted row. The single row must have been inserted rather than updated. Otherwise `oid` is zero.

If the `INSERT` command contains a `RETURNING` clause, the result will be similar to that of a `SELECT` statement containing the columns and values defined in the `RETURNING` list, computed over the row(s) inserted or updated by the command.

Examples

Insert a single row into table `films`:

```
INSERT INTO films VALUES ('UA502', 'Bananas', 105, '1971-07-13', 'Comedy', '82 minutes');
```

In this example, the `len` column is omitted and therefore it will have the default value:

```
INSERT INTO films (code, title, did, date_prod, kind) VALUES ('T_601', 'Yojimbo', 106, '1961-06-16', 'Drama');
```

This example uses the `DEFAULT` clause for the date columns rather than specifying a value:

```
INSERT INTO films VALUES ('UA502', 'Bananas', 105, DEFAULT, 'Comedy', '82 minutes');
INSERT INTO films (code, title, did, date_prod, kind) VALUES ('T_601', 'Yojimbo', 106, DEFAULT, 'Drama');
```

To insert a row consisting entirely of default values:

```
INSERT INTO films DEFAULT VALUES;
```

To insert multiple rows using the multirow `VALUES` syntax:

```
INSERT INTO films (code, title, did, date_prod, kind) VALUES ('B6717', 'Tampopo', 110, '1985-02-10', 'Comedy'),
 ('HG120', 'The Dinner Game', 140, DEFAULT, 'Comedy');
```

This example inserts some rows into table `films` from a table `tmp_films` with the same column layout as `films`:

```
INSERT INTO films SELECT * FROM tmp_films WHERE date_prod < '2004-05-07';
```
This example inserts into array columns:

```
-- Create an empty 3x3 gameboard for noughts-and-crosses
INSERT INTO tictactoe (game, board[1:3][1:3])
VALUES (1, '{{"","",""},{"","",""},{"","",""}}');
-- The subscripts in the above example aren't really needed
INSERT INTO tictactoe (game, board)
VALUES (2, '{{X,"",""},{","O",""},{","X",""}}');
```

Insert a single row into table `distributors`, returning the sequence number generated by the `DEFAULT` clause:

```
INSERT INTO distributors (did, dname) VALUES (DEFAULT, 'XYZ Widgets')
RETURNING did;
```

Increment the sales count of the salesperson who manages the account for Acme Corporation, and record the whole updated row along with current time in a log table:

```
WITH upd AS (
  UPDATE employees SET sales_count = sales_count + 1 WHERE id =
  (SELECT sales_person FROM accounts WHERE name = 'Acme Corporation')
  RETURNING *
)
INSERT INTO employees_log SELECT *, current_timestamp FROM upd;
```

Insert or update new distributors as appropriate. Assumes a unique index has been defined that constrains values appearing in the `did` column. Note that the special `excluded` table is used to reference values originally proposed for insertion:

```
INSERT INTO distributors (did, dname)
VALUES (5, 'Gizmo Transglobal'), (6, 'Associated Computing, Inc')
ON CONFLICT (did) DO UPDATE SET dname = EXCLUDED.dname;
```

Insert a distributor, or do nothing for rows proposed for insertion when an existing, excluded row (a row with a matching constrained column or columns after before row insert triggers fire) exists. Example assumes a unique index has been defined that constrains values appearing in the `did` column:

```
INSERT INTO distributors (did, dname) VALUES (7, 'Redline GmbH')
ON CONFLICT (did) DO NOTHING;
```

Insert or update new distributors as appropriate. Example assumes a unique index has been defined that constrains values appearing in the `did` column. `WHERE` clause is used to limit the rows actually updated (any existing row not updated will still be locked, though):

```
-- Don't update existing distributors based in a certain ZIP code
INSERT INTO distributors AS d (did, dname) VALUES (8, 'Anvil Distribution')
ON CONFLICT (did) DO UPDATE
SET dname = EXCLUDED.dname || ' (formerly ' || d.dname || ')'
WHERE d.zipcode <> '21201';
```
-- Name a constraint directly in the statement (uses associated
-- index to arbitrate taking the DO NOTHING action)
INSERT INTO distributors (did, dname) VALUES (9, 'Antwerp Design')
  ON CONFLICT ON CONSTRAINT distributors_pkey DO NOTHING;

Insert new distributor if possible; otherwise DO NOTHING. Example assumes a unique index has been
defined that constrains values appearing in the did column on a subset of rows where the is_active
Boolean column evaluates to true:

-- This statement could infer a partial unique index on "did"
-- with a predicate of "WHERE is_active", but it could also
-- just use a regular unique constraint on "did"
INSERT INTO distributors (did, dname) VALUES (10, 'Conrad International')
  ON CONFLICT (did) WHERE is_active DO NOTHING;

Compatibility

INSERT conforms to the SQL standard, except that the RETURNING clause is a PostgreSQL extension,
as is the ability to use WITH with INSERT. and the ability to specify an alternative action with
ON CONFLICT. Also, the case in which a column name list is omitted, but not all the columns are filled
from the VALUES clause or query, is disallowed by the standard.

Possible limitations of the query clause are documented under SELECT.
LISTEN

Name
LISTEN — listen for a notification

Synopsis
LISTEN channel

Description
LISTEN registers the current session as a listener on the notification channel named channel. If the current session is already registered as a listener for this notification channel, nothing is done.

Whenever the command NOTIFY channel is invoked, either by this session or another one connected to the same database, all the sessions currently listening on that notification channel are notified, and each will in turn notify its connected client application.

A session can be unregistered for a given notification channel with the UNLISTEN command. A session’s listen registrations are automatically cleared when the session ends.

The method a client application must use to detect notification events depends on which PostgreSQL application programming interface it uses. With the libpq library, the application issues LISTEN as an ordinary SQL command, and then must periodically call the function PQnotifies to find out whether any notification events have been received. Other interfaces such as libpgcrl provide higher-level methods for handling notify events; indeed, with libpgcrl the application programmer should not even issue LISTEN or UNLISTEN directly. See the documentation for the interface you are using for more details.

NOTIFY contains a more extensive discussion of the use of LISTEN and NOTIFY.

Parameters

channel

Name of a notification channel (any identifier).

Notes
LISTEN takes effect at transaction commit. If LISTEN or UNLISTEN is executed within a transaction that later rolls back, the set of notification channels being listened to is unchanged.

A transaction that has executed LISTEN cannot be prepared for two-phase commit.
Examples
Configure and execute a listen/notify sequence from psql:

LISTEN virtual;
NOTIFY virtual;
Asynchronous notification "virtual" received from server process with PID 8448.

Compatibility
There is no LISTEN statement in the SQL standard.

See Also
NOTIFY, UNLISTEN
LOAD

Name
LOAD — load a shared library file

Synopsis
LOAD 'filename'

Description
This command loads a shared library file into the PostgreSQL server's address space. If the file has been loaded already, the command does nothing. Shared library files that contain C functions are automatically loaded whenever one of their functions is called. Therefore, an explicit LOAD is usually only needed to load a library that modifies the server's behavior through "hooks" rather than providing a set of functions.

The file name is specified in the same way as for shared library names in CREATE FUNCTION; in particular, one can rely on a search path and automatic addition of the system's standard shared library file name extension. See Section 36.9 for more information on this topic.

Non-superusers can only apply LOAD to library files located in $libdir/plugins/ — the specified filename must begin with exactly that string. (It is the database administrator's responsibility to ensure that only "safe" libraries are installed there.)

Compatibility
LOAD is a PostgreSQL extension.

See Also
CREATE FUNCTION
LOCK

Name
LOCK — lock a table

Synopsis
LOCK [ TABLE ] [ ONLY ] name [ * ] [, ...] [ IN lockmode MODE ] [ NOWAIT ]

where lockmode is one of:

ACCESS SHARE | ROW SHARE | ROW EXCLUSIVE | SHARE UPDATE EXCLUSIVE |
| SHARE | SHARE ROW EXCLUSIVE | EXCLUSIVE | ACCESS EXCLUSIVE

Description
LOCK TABLE obtains a table-level lock, waiting if necessary for any conflicting locks to be released. If NOWAIT is specified, LOCK TABLE does not wait to acquire the desired lock: if it cannot be acquired immediately, the command is aborted and an error is emitted. Once obtained, the lock is held for the remainder of the current transaction. (There is no UNLOCK TABLE command; locks are always released at transaction end.)

When acquiring locks automatically for commands that reference tables, PostgreSQL always uses the least restrictive lock mode possible. LOCK TABLE provides for cases when you might need more restrictive locking. For example, suppose an application runs a transaction at the READ COMMITTED isolation level and needs to ensure that data in a table remains stable for the duration of the transaction. To achieve this you could obtain SHARE lock mode over the table before querying. This will prevent concurrent data changes and ensure subsequent reads of the table see a stable view of committed data, because SHARE lock mode conflicts with the ROW EXCLUSIVE lock acquired by writers, and your LOCK TABLE name IN SHARE MODE statement will wait until any concurrent holders of ROW EXCLUSIVE mode locks commit or roll back. Thus, once you obtain the lock, there are no uncommitted writes outstanding; furthermore none can begin until you release the lock.

To achieve a similar effect when running a transaction at the REPEATABLE READ or SERIALIZABLE isolation level, you have to execute the LOCK TABLE statement before executing any SELECT or data modification statement. A REPEATABLE READ or SERIALIZABLE transaction’s view of data will be frozen when its first SELECT or data modification statement begins. A LOCK TABLE later in the transaction will still prevent concurrent writes — but it won’t ensure that what the transaction reads corresponds to the latest committed values.

If a transaction of this sort is going to change the data in the table, then it should use SHARE ROW EXCLUSIVE lock mode instead of SHARE mode. This ensures that only one transaction of this type runs at a time. Without this, a deadlock is possible: two transactions might both acquire SHARE mode, and then be unable to also acquire ROW EXCLUSIVE mode to actually perform their updates. (Note that a transaction’s own locks never conflict, so a transaction can acquire ROW EXCLUSIVE mode when it holds SHARE mode — but not if anyone else holds SHARE mode.) To avoid deadlocks, make sure all transactions acquire locks on the same objects in the same order, and if multiple lock modes are involved for a single object, then transactions should always acquire the most restrictive mode first.

More information about the lock modes and locking strategies can be found in Section 13.3.
Parameters

name

The name (optionally schema-qualified) of an existing table to lock. If ONLY is specified before the table name, only that table is locked. If ONLY is not specified, the table and all its descendant tables (if any) are locked. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

The command LOCK TABLE a, b; is equivalent to LOCK TABLE a; LOCK TABLE b; The tables are locked one-by-one in the order specified in the LOCK TABLE command.

lockmode

The lock mode specifies which locks this lock conflicts with. Lock modes are described in Section 13.3.

If no lock mode is specified, then ACCESS EXCLUSIVE, the most restrictive mode, is used.

NOWAIT

Specifies that LOCK TABLE should not wait for any conflicting locks to be released: if the specified lock(s) cannot be acquired immediately without waiting, the transaction is aborted.

Notes

LOCK TABLE ... IN ACCESS SHARE MODE requires SELECT privileges on the target table. LOCK TABLE ... IN ROW EXCLUSIVE MODE requires INSERT, UPDATE, DELETE, or TRUNCATE privileges on the target table. All other forms of LOCK require table-level UPDATE, DELETE, or TRUNCATE privileges.

LOCK TABLE is useless outside a transaction block: the lock would remain held only to the completion of the statement. Therefore PostgreSQL reports an error if LOCK is used outside a transaction block. Use BEGIN and COMMIT (or ROLLBACK) to define a transaction block.

LOCK TABLE only deals with table-level locks, and so the mode names involving ROW are all misnomers. These mode names should generally be read as indicating the intention of the user to acquire row-level locks within the locked table. Also, ROW EXCLUSIVE mode is a shareable table lock. Keep in mind that all the lock modes have identical semantics so far as LOCK TABLE is concerned, differing only in the rules about which modes conflict with which. For information on how to acquire an actual row-level lock, see Section 13.3.2 and the The Locking Clause in the SELECT reference documentation.

Examples

Obtain a SHARE lock on a primary key table when going to perform inserts into a foreign key table:

BEGIN WORK;
LOCK TABLE films IN SHARE MODE;
SELECT id FROM films
    WHERE name = 'Star Wars: Episode I - The Phantom Menace';
-- Do ROLLBACK if record was not returned
INSERT INTO films_user_comments VALUES
    (_id_, 'GREAT! I was waiting for it for so long!');
COMMIT WORK;
Take a \texttt{SHARE ROW EXCLUSIVE} lock on a primary key table when going to perform a delete operation:

\begin{verbatim}
BEGIN WORK;
LOCK TABLE films IN SHARE ROW EXCLUSIVE MODE;
DELETE FROM films_user_comments WHERE id IN
   (SELECT id FROM films WHERE rating < 5);
DELETE FROM films WHERE rating < 5;
COMMIT WORK;
\end{verbatim}

\section*{Compatibility}

There is no \texttt{LOCK TABLE} in the SQL standard, which instead uses \texttt{SET TRANSACTION} to specify concurrency levels on transactions. PostgreSQL supports that too; see \texttt{SET TRANSACTION} for details.

Except for \texttt{ACCESS SHARE}, \texttt{ACCESS EXCLUSIVE}, and \texttt{SHARE UPDATE EXCLUSIVE} lock modes, the PostgreSQL lock modes and the \texttt{LOCK TABLE} syntax are compatible with those present in Oracle.
MOVE

Name
MOVE — position a cursor

Synopsis
MOVE [ direction [ FROM | IN ] ] cursor_name

where direction can be empty or one of:

   NEXT
   PRIOR
   FIRST
   LAST
   ABSOLUTE count
   RELATIVE count
   count
   ALL
   FORWARD
   FORWARD count
   FORWARD ALL
   BACKWARD
   BACKWARD count
   BACKWARD ALL

Description
MOVE repositions a cursor without retrieving any data. MOVE works exactly like the FETCH command, except it only positions the cursor and does not return rows.

The parameters for the MOVE command are identical to those of the FETCH command; refer to FETCH for details on syntax and usage.

Outputs
On successful completion, a MOVE command returns a command tag of the form

MOVE count

The count is the number of rows that a FETCH command with the same parameters would have returned (possibly zero).

Examples
BEGIN WORK;
DECLARE liahona CURSOR FOR SELECT * FROM films;

   -- Skip the first 5 rows:
MOVE FORWARD 5 IN liahona;
MOVE 5

-- Fetch the 6th row from the cursor liahona:
FETCH 1 FROM liahona;

<table>
<thead>
<tr>
<th>code</th>
<th>title</th>
<th>did</th>
<th>date_prod</th>
<th>kind</th>
<th>len</th>
</tr>
</thead>
<tbody>
<tr>
<td>P_303</td>
<td>48 Hrs</td>
<td>103</td>
<td>1982-10-22</td>
<td>Action</td>
<td>01:37</td>
</tr>
</tbody>
</table>

(1 row)

-- Close the cursor liahona and end the transaction:
CLOSE liahona;
COMMIT WORK;

Compatibility

There is no MOVE statement in the SQL standard.

See Also

CLOSE, DECLARE, FETCH
NOTIFY

Name

NOTIFY — generate a notification

Synopsis

NOTIFY channel [ , payload ]

Description

The NOTIFY command sends a notification event together with an optional “payload” string to each client application that has previously executed LISTEN channel for the specified channel name in the current database. Notifications are visible to all users.

NOTIFY provides a simple interprocess communication mechanism for a collection of processes accessing the same PostgreSQL database. A payload string can be sent along with the notification, and higher-level mechanisms for passing structured data can be built by using tables in the database to pass additional data from notifier to listener(s).

The information passed to the client for a notification event includes the notification channel name, the notifying session’s server process PID, and the payload string, which is an empty string if it has not been specified.

It is up to the database designer to define the channel names that will be used in a given database and what each one means. Commonly, the channel name is the same as the name of some table in the database, and the notify event essentially means, “I changed this table, take a look at it to see what’s new”. But no such association is enforced by the NOTIFY and LISTEN commands. For example, a database designer could use several different channel names to signal different sorts of changes to a single table. Alternatively, the payload string could be used to differentiate various cases.

When NOTIFY is used to signal the occurrence of changes to a particular table, a useful programming technique is to put the NOTIFY in a statement trigger that is triggered by table updates. In this way, notification happens automatically when the table is changed, and the application programmer cannot accidentally forget to do it.

NOTIFY interacts with SQL transactions in some important ways. Firstly, if a NOTIFY is executed inside a transaction, the notify events are not delivered until and unless the transaction is committed. This is appropriate, since if the transaction is aborted, all the commands within it have had no effect, including NOTIFY. But it can be disconcerting if one is expecting the notification events to be delivered immediately. Secondly, if a listening session receives a notification signal while it is within a transaction, the notification event will not be delivered to its connected client until just after the transaction is completed (either committed or aborted). Again, the reasoning is that if a notification were delivered within a transaction that was later aborted, one would want the notification to be undone somehow — but the server cannot “take back” a notification once it has sent it to the client. So notification events are only delivered between transactions. The upshot of this is that applications using NOTIFY for real-time signaling should try to keep their transactions short.

If the same channel name is signaled multiple times from the same transaction with identical payload strings, the database server can decide to deliver a single notification only. On the other hand, notifications with distinct payload strings will always be delivered as distinct notifications. Similarly,
notifications from different transactions will never get folded into one notification. Except for dropping later instances of duplicate notifications, NOTIFY guarantees that notifications from the same transaction get delivered in the order they were sent. It is also guaranteed that messages from different transactions are delivered in the order in which the transactions committed.

It is common for a client that executes NOTIFY to be listening on the same notification channel itself. In that case it will get back a notification event, just like all the other listening sessions. Depending on the application logic, this could result in useless work, for example, reading a database table to find the same updates that that session just wrote out. It is possible to avoid such extra work by noticing whether the notifying session’s server process PID (supplied in the notification event message) is the same as one’s own session’s PID (available from libpq). When they are the same, the notification event is one’s own work bouncing back, and can be ignored.

**Parameters**

*channel*

Name of the notification channel to be signaled (any identifier).

*payload*

The “payload” string to be communicated along with the notification. This must be specified as a simple string literal. In the default configuration it must be shorter than 8000 bytes. (If binary data or large amounts of information need to be communicated, it’s best to put it in a database table and send the key of the record.)

**Notes**

There is a queue that holds notifications that have been sent but not yet processed by all listening sessions. If this queue becomes full, transactions calling NOTIFY will fail at commit. The queue is quite large (8GB in a standard installation) and should be sufficiently sized for almost every use case. However, no cleanup can take place if a session executes LISTEN and then enters a transaction for a very long time. Once the queue is half full you will see warnings in the log file pointing you to the session that is preventing cleanup. In this case you should make sure that this session ends its current transaction so that cleanup can proceed.

The function `pg_notification_queue_usage` returns the fraction of the queue that is currently occupied by pending notifications. See Section 9.25 for more information.

A transaction that has executed NOTIFY cannot be prepared for two-phase commit.

**pg_notify**

To send a notification you can also use the function `pg_notify(text, text)`. The function takes the channel name as the first argument and the payload as the second. The function is much easier to use than the NOTIFY command if you need to work with non-constant channel names and payloads.
Examples

Configure and execute a listen/notify sequence from psql:

```
LISTEN virtual;
NOTIFY virtual;
```

Asynchronous notification "virtual" received from server process with PID 8448.
```
NOTIFY virtual, 'This is the payload';
```
Asynchronous notification "virtual" with payload "This is the payload" received from server process with PID 8448.
```
LISTEN foo;
SELECT pg_notify('fo'||'o', 'pay'||'load');
```
Asynchronous notification "foo" with payload "payload" received from server process with PID 14728.

Compatibility

There is no NOTIFY statement in the SQL standard.

See Also

LISTEN, UNLISTEN
**PREPARE**

**Name**

PREPARE — prepare a statement for execution

**Synopsis**

PREPARE name [ ( data_type [, ...] ) ] AS statement

**Description**

PREPARE creates a prepared statement. A prepared statement is a server-side object that can be used to optimize performance. When the PREPARE statement is executed, the specified statement is parsed, analyzed, and rewritten. When an EXECUTE command is subsequently issued, the prepared statement is planned and executed. This division of labor avoids repetitive parse analysis work, while allowing the execution plan to depend on the specific parameter values supplied.

Prepared statements can take parameters: values that are substituted into the statement when it is executed. When creating the prepared statement, refer to parameters by position, using $1, $2, etc. A corresponding list of parameter data types can optionally be specified. When a parameter’s data type is not specified or is declared as unknown, the type is inferred from the context in which the parameter is first used (if possible). When executing the statement, specify the actual values for these parameters in the EXECUTE statement. Refer to EXECUTE for more information about that.

Prepared statements only last for the duration of the current database session. When the session ends, the prepared statement is forgotten, so it must be recreated before being used again. This also means that a single prepared statement cannot be used by multiple simultaneous database clients; however, each client can create their own prepared statement to use. Prepared statements can be manually cleaned up using the DEALLOCATE command.

Prepared statements potentially have the largest performance advantage when a single session is being used to execute a large number of similar statements. The performance difference will be particularly significant if the statements are complex to plan or rewrite, e.g. if the query involves a join of many tables or requires the application of several rules. If the statement is relatively simple to plan and rewrite but relatively expensive to execute, the performance advantage of prepared statements will be less noticeable.

**Parameters**

**name**

An arbitrary name given to this particular prepared statement. It must be unique within a single session and is subsequently used to execute or deallocate a previously prepared statement.

**data_type**

The data type of a parameter to the prepared statement. If the data type of a particular parameter is unspecified or is specified as unknown, it will be inferred from the context in which the parameter is first used. To refer to the parameters in the prepared statement itself, use $1, $2, etc.
PREPARE

statement

Any SELECT, INSERT, UPDATE, DELETE, or VALUES statement.

Notes

Prepared statements can use generic plans rather than re-planning with each set of supplied EXECUTE values. This occurs immediately for prepared statements with no parameters; otherwise it occurs only after five or more executions produce plans whose estimated cost average (including planning overhead) is more expensive than the generic plan cost estimate. Once a generic plan is chosen, it is used for the remaining lifetime of the prepared statement. Using EXECUTE values which are rare in columns with many duplicates can generate custom plans that are so much cheaper than the generic plan, even after adding planning overhead, that the generic plan might never be used.

A generic plan assumes that each value supplied to EXECUTE is one of the column’s distinct values and that column values are uniformly distributed. For example, if statistics record three distinct column values, a generic plan assumes a column equality comparison will match 33% of processed rows. Column statistics also allow generic plans to accurately compute the selectivity of unique columns. Comparisons on non-uniformly-distributed columns and specification of non-existent values affects the average plan cost, and hence if and when a generic plan is chosen.

To examine the query plan PostgreSQL is using for a prepared statement, use EXPLAIN, e.g. EXPLAIN EXECUTE. If a generic plan is in use, it will contain parameter symbols $n, while a custom plan will have the supplied parameter values substituted into it. The row estimates in the generic plan reflect the selectivity computed for the parameters.

For more information on query planning and the statistics collected by PostgreSQL for that purpose, see the ANALYZE documentation.

Although the main point of a prepared statement is to avoid repeated parse analysis and planning of the statement, PostgreSQL will force re-analysis and re-planning of the statement before using it whenever database objects used in the statement have undergone definitional (DDL) changes since the previous use of the prepared statement. Also, if the value of search_path changes from one use to the next, the statement will be re-parsed using the new search_path. (This latter behavior is new as of PostgreSQL 9.3.) These rules make use of a prepared statement semantically almost equivalent to re-submitting the same query text over and over, but with a performance benefit if no object definitions are changed, especially if the best plan remains the same across uses. An example of a case where the semantic equivalence is not perfect is that if the statement refers to a table by an unqualified name, and then a new table of the same name is created in a schema appearing earlier in the search_path, no automatic re-parse will occur since no object used in the statement changed. However, if some other change forces a re-parse, the new table will be referenced in subsequent uses.

You can see all prepared statements available in the session by querying the pg_prepared_statements system view.

Examples

Create a prepared statement for an INSERT statement, and then execute it:

PREPARE fooplan (int, text, bool, numeric) AS
    INSERT INTO foo VALUES($1, $2, $3, $4);
EXECUTE fooplan(1, 'Hunter Valley', 't', 200.00);
Prepare

Create a prepared statement for a `SELECT` statement, and then execute it:

```sql
PREPARE usrrptplan (int) AS
    SELECT * FROM users u, logs l WHERE u.usrid=$1 AND u.usrid=l.usrid
    AND l.date = $2;
EXECUTE usrrptplan(1, current_date);
```

Note that the data type of the second parameter is not specified, so it is inferred from the context in which $2 is used.

Compatibility

The SQL standard includes a `PREPARE` statement, but it is only for use in embedded SQL. This version of the `PREPARE` statement also uses a somewhat different syntax.

See Also

`DEALLOCATE`, `EXECUTE`
PREPARE TRANSACTION

Name

PREPARE TRANSACTION — prepare the current transaction for two-phase commit

Synopsis

PREPARE TRANSACTION transaction_id

Description

PREPARE TRANSACTION prepares the current transaction for two-phase commit. After this command, the transaction is no longer associated with the current session; instead, its state is fully stored on disk, and there is a very high probability that it can be committed successfully, even if a database crash occurs before the commit is requested.

Once prepared, a transaction can later be committed or rolled back with COMMIT PREPARED or ROLLBACK PREPARED, respectively. Those commands can be issued from any session, not only the one that executed the original transaction.

From the point of view of the issuing session, PREPARE TRANSACTION is not unlike a ROLLBACK command: after executing it, there is no active current transaction, and the effects of the prepared transaction are no longer visible. (The effects will become visible again if the transaction is committed.)

If the PREPARE TRANSACTION command fails for any reason, it becomes a ROLLBACK: the current transaction is canceled.

Parameters

transaction_id

An arbitrary identifier that later identifies this transaction for COMMIT PREPARED or ROLLBACK PREPARED. The identifier must be written as a string literal, and must be less than 200 bytes long. It must not be the same as the identifier used for any currently prepared transaction.

Notes

PREPARE TRANSACTION is not intended for use in applications or interactive sessions. Its purpose is to allow an external transaction manager to perform atomic global transactions across multiple databases or other transactional resources. Unless you’re writing a transaction manager, you probably shouldn’t be using PREPARE TRANSACTION.

This command must be used inside a transaction block. Use BEGIN to start one.

It is not currently allowed to PREPARE a transaction that has executed any operations involving temporary tables, created any cursors WITH HOLD, or executed LISTEN, UNLISTEN, or NOTIFY. Those features are too tightly tied to the current session to be useful in a transaction to be prepared.
If the transaction modified any run-time parameters with `SET` (without the `LOCAL` option), those effects persist after `PREPARE TRANSACTION`, and will not be affected by any later `COMMIT PREPARED` or `ROLLBACK PREPARED`. Thus, in this one respect `PREPARE TRANSACTION` acts more like `COMMIT` than `ROLLBACK`.

All currently available prepared transactions are listed in the `pg_prepared_xacts` system view.

Caution

It is unwise to leave transactions in the prepared state for a long time. This will interfere with the ability of `VACUUM` to reclaim storage, and in extreme cases could cause the database to shut down to prevent transaction ID wraparound (see Section 24.1.5). Keep in mind also that the transaction continues to hold whatever locks it held. The intended usage of the feature is that a prepared transaction will normally be committed or rolled back as soon as an external transaction manager has verified that other databases are also prepared to commit.

If you have not set up an external transaction manager to track prepared transactions and ensure they get closed out promptly, it is best to keep the prepared-transaction feature disabled by setting `max_prepared_transactions` to zero. This will prevent accidental creation of prepared transactions that might then be forgotten and eventually cause problems.

Examples

Prepare the current transaction for two-phase commit, using `foobar` as the transaction identifier:

```
PREPARE TRANSACTION 'foobar';
```

Compatibility

`PREPARE TRANSACTION` is a PostgreSQL extension. It is intended for use by external transaction management systems, some of which are covered by standards (such as X/Open XA), but the SQL side of those systems is not standardized.

See Also

`COMMIT PREPARED`, `ROLLBACK PREPARED`
REASSIGN OWNED

Name

REASSIGN OWNED — change the ownership of database objects owned by a database role

Synopsis

REASSIGN OWNED BY { old_role | CURRENT_USER | SESSION_USER } [, ...]
TO { new_role | CURRENT_USER | SESSION_USER }

Description

REASSIGN OWNED instructs the system to change the ownership of database objects owned by any of the old_roles to new_role.

Parameters

old_role

The name of a role. The ownership of all the objects within the current database, and of all shared objects (databases, tablespaces), owned by this role will be reassigned to new_role.

new_role

The name of the role that will be made the new owner of the affected objects.

Notes

REASSIGN OWNED is often used to prepare for the removal of one or more roles. Because REASSIGN OWNED does not affect objects within other databases, it is usually necessary to execute this command in each database that contains objects owned by a role that is to be removed.

REASSIGN OWNED requires privileges on both the source role(s) and the target role.

The DROP OWNED command is an alternative that simply drops all the database objects owned by one or more roles.

The REASSIGN OWNED command does not affect any privileges granted to the old_roles on objects that are not owned by them. Likewise, it does not affect default privileges created with ALTER DEFAULT PRIVILEGES. Use DROP OWNED to revoke such privileges.

See Section 21.4 for more discussion.

Compatibility

The REASSIGN OWNED command is a PostgreSQL extension.
See Also

DROP OWNED, DROP ROLE, ALTER DATABASE
REFRESH MATERIALIZED VIEW

Name

REFRESH MATERIALIZED VIEW — replace the contents of a materialized view

Synopsis

REFRESH MATERIALIZED VIEW [ CONCURRENTLY ] name
[ WITH [ NO ] DATA ]

Description

REFRESH MATERIALIZED VIEW completely replaces the contents of a materialized view. To execute this command you must be the owner of the materialized view. The old contents are discarded. If WITH DATA is specified (or defaults) the backing query is executed to provide the new data, and the materialized view is left in a scannable state. If WITH NO DATA is specified no new data is generated and the materialized view is left in an unscannable state.

CONCURRENTLY and WITH NO DATA may not be specified together.

Parameters

CONCURRENTLY

Refresh the materialized view without locking out concurrent selects on the materialized view. Without this option a refresh which affects a lot of rows will tend to use fewer resources and complete more quickly, but could block other connections which are trying to read from the materialized view. This option may be faster in cases where a small number of rows are affected.

This option is only allowed if there is at least one UNIQUE index on the materialized view which uses only column names and includes all rows; that is, it must not be an expression index or include a WHERE clause.

This option may not be used when the materialized view is not already populated.

Even with this option only one REFRESH at a time may run against any one materialized view.

name

The name (optionally schema-qualified) of the materialized view to refresh.

Notes

While the default index for future CLUSTER operations is retained, REFRESH MATERIALIZED VIEW does not order the generated rows based on this property. If you want the data to be ordered upon generation, you must use an ORDER BY clause in the backing query.
Examples

This command will replace the contents of the materialized view called order_summary using the query from the materialized view’s definition, and leave it in a scannable state:

REFRESH MATERIALIZED VIEW order_summary;

This command will free storage associated with the materialized view annual_statistics_basis and leave it in an unscannable state:

REFRESH MATERIALIZED VIEW annual_statistics_basis WITH NO DATA;

Compatibility

REFRESH MATERIALIZED VIEW is a PostgreSQL extension.

See Also

CREATE MATERIALIZED VIEW, ALTER MATERIALIZED VIEW, DROP MATERIALIZED VIEW
**REINDEX**

**Name**

REINDEX — rebuild indexes

**Synopsis**

REINDEX [ (VERBOSE) ] [ INDEX | TABLE | SCHEMA | DATABASE | SYSTEM ] name

**Description**

REINDEX rebuilds an index using the data stored in the index’s table, replacing the old copy of the index. There are several scenarios in which to use REINDEX:

- An index has become corrupted, and no longer contains valid data. Although in theory this should never happen, in practice indexes can become corrupted due to software bugs or hardware failures. REINDEX provides a recovery method.
- An index has become “bloated”, that is it contains many empty or nearly-empty pages. This can occur with B-tree indexes in PostgreSQL under certain uncommon access patterns. REINDEX provides a way to reduce the space consumption of the index by writing a new version of the index without the dead pages. See Section 24.2 for more information.
- You have altered a storage parameter (such as fillfactor) for an index, and wish to ensure that the change has taken full effect.
- An index build with the CONCURRENTLY option failed, leaving an “invalid” index. Such indexes are useless but it can be convenient to use REINDEX to rebuild them. Note that REINDEX will not perform a concurrent build. To build the index without interfering with production you should drop the index and reissue the CREATE INDEX CONCURRENTLY command.

**Parameters**

INDEX

Recreate the specified index.

TABLE

Recreate all indexes of the specified table. If the table has a secondary “TOAST” table, that is reindexed as well.

SCHEMA

Recreate all indexes of the specified schema. If a table of this schema has a secondary “TOAST” table, that is reindexed as well. Indexes on shared system catalogs are also processed. This form of REINDEX cannot be executed inside a transaction block.
REINDEX

DATABASE

Recreate all indexes within the current database. Indexes on shared system catalogs are also processed. This form of REINDEX cannot be executed inside a transaction block.

SYSTEM

Recreate all indexes on system catalogs within the current database. Indexes on shared system catalogs are included. Indexes on user tables are not processed. This form of REINDEX cannot be executed inside a transaction block.

name

The name of the specific index, table, or database to be reindexed. Index and table names can be schema-qualified. Presently, REINDEX DATABASE and REINDEX SYSTEM can only reindex the current database, so their parameter must match the current database’s name.

VERBOSE

Prints a progress report as each index is reindexed.

Notes

If you suspect corruption of an index on a user table, you can simply rebuild that index, or all indexes on the table, using REINDEX INDEX or REINDEX TABLE.

Things are more difficult if you need to recover from corruption of an index on a system table. In this case it’s important for the system to not have used any of the suspect indexes itself. (Indeed, in this sort of scenario you might find that server processes are crashing immediately at start-up, due to reliance on the corrupted indexes.) To recover safely, the server must be started with the -P option, which prevents it from using indexes for system catalog lookups.

One way to do this is to shut down the server and start a single-user PostgreSQL server with the -P option included on its command line. Then, REINDEX DATABASE, REINDEX SYSTEM, REINDEX TABLE, or REINDEX INDEX can be issued, depending on how much you want to reconstruct. If in doubt, use REINDEX SYSTEM to select reconstruction of all system indexes in the database. Then quit the single-user server session and restart the regular server. See the postgres reference page for more information about how to interact with the single-user server interface.

Alternatively, a regular server session can be started with -P included in its command line options. The method for doing this varies across clients, but in all libpq-based clients, it is possible to set the PGOPTIONS environment variable to -P before starting the client. Note that while this method does not require locking out other clients, it might still be wise to prevent other users from connecting to the damaged database until repairs have been completed.

REINDEX is similar to a drop and recreate of the index in that the index contents are rebuilt from scratch. However, the locking considerations are rather different. REINDEX locks out writes but not reads of the index’s parent table. It also takes an exclusive lock on the specific index being processed, which will block reads that attempt to use that index. In contrast, DROP INDEX momentarily takes an exclusive lock on the parent table, blocking both writes and reads. The subsequent CREATE INDEX locks out writes but not reads; since the index is not there, no read will attempt to use it, meaning that there will be no blocking but reads might be forced into expensive sequential scans.

Reindexing a single index or table requires being the owner of that index or table. Reindexing a database requires being the owner of the database (note that the owner can therefore rebuild indexes of tables owned by other users). Of course, superusers can always reindex anything.
Examples

Rebuild a single index:

REINDEX INDEX my_index;

Rebuild all the indexes on the table my_table:

REINDEX TABLE my_table;

Rebuild all indexes in a particular database, without trusting the system indexes to be valid already:

$ export PGOPTIONS="-p"
$ psql broken_db
...
broken_db=> REINDEX DATABASE broken_db;
broken_db=> \q

Compatibility

There is no REINDEX command in the SQL standard.
RELEASE SAVEPOINT

Name
RELEASE SAVEPOINT — destroy a previously defined savepoint

Synopsis
RELEASE [ SAVEPOINT ] savepoint_name

Description
RELEASE SAVEPOINT destroys a savepoint previously defined in the current transaction.
Destroying a savepoint makes it unavailable as a rollback point, but it has no other user visible behavior. It does not undo the effects of commands executed after the savepoint was established. (To do that, see ROLLBACK TO SAVEPOINT.) Destroying a savepoint when it is no longer needed allows the system to reclaim some resources earlier than transaction end.
RELEASE SAVEPOINT also destroys all savepoints that were established after the named savepoint was established.

Parameters
savepoint_name
The name of the savepoint to destroy.

Notes
Specifying a savepoint name that was not previously defined is an error.
It is not possible to release a savepoint when the transaction is in an aborted state.
If multiple savepoints have the same name, only the one that was most recently defined is released.

Examples
To establish and later destroy a savepoint:

BEGIN;
    INSERT INTO table1 VALUES (3);
    SAVEPOINT my_savepoint;
    INSERT INTO table1 VALUES (4);
    RELEASE SAVEPOINT my_savepoint;
COMMIT;

The above transaction will insert both 3 and 4.
Compatibility

This command conforms to the SQL standard. The standard specifies that the key word `SAVEPOINT` is mandatory, but PostgreSQL allows it to be omitted.

See Also

BEGIN, COMMIT, ROLLBACK, ROLLBACK TO SAVEPOINT, SAVEPOINT
RESET

Name

RESET — restore the value of a run-time parameter to the default value

Synopsis

```
RESET configuration_parameter
RESET ALL
```

Description

RESET restores run-time parameters to their default values. RESET is an alternative spelling for

```
SET configuration_parameter TO DEFAULT
```

Refer to SET for details.

The default value is defined as the value that the parameter would have had, if no SET had ever been
issued for it in the current session. The actual source of this value might be a compiled-in default, the
configuration file, command-line options, or per-database or per-user default settings. This is subtly
different from defining it as “the value that the parameter had at session start”, because if the value
came from the configuration file, it will be reset to whatever is specified by the configuration file now.
See Chapter 19 for details.

The transactional behavior of RESET is the same as SET: its effects will be undone by transaction
rollback.

Parameters

configuration_parameter

Name of a settable run-time parameter. Available parameters are documented in Chapter 19 and
on the SET reference page.

ALL

Resets all settable run-time parameters to default values.

Examples

Set the timezone configuration variable to its default value:

```
RESET timezone;
```
Compatibility

RESET is a PostgreSQL extension.

See Also

SET, SHOW
REVOKE

Name

REVOKE — remove access privileges

Synopsis

REVOKE [ GRANT OPTION FOR ]

{ { SELECT | INSERT | UPDATE | DELETE | TRUNCATE | REFERENCES | TRIGGER }

[, ...] | ALL [ PRIVILEGES ] }
ON { [ TABLE ] table_name [, ...]

| ALL TABLES IN SCHEMA schema_name [, ...] }
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]

{ { SELECT | INSERT | UPDATE | REFERENCES }

( column_name [, ...] )

[, ...] | ALL [ PRIVILEGES ] ( column_name [, ...] ) }
ON [ TABLE ] table_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]

{ { USAGE | SELECT | UPDATE }

[, ...] | ALL [ PRIVILEGES ] }
ON { SEQUENCE sequence_name [, ...]

| ALL SEQUENCES IN SCHEMA schema_name [, ...] }
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]

{ { CREATE | CONNECT | TEMPORARY | TEMP }

[, ...] | ALL [ PRIVILEGES ] }
ON DATABASE database_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]

{ USAGE | ALL [ PRIVILEGES ] }
ON DOMAIN domain_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]

{ USAGE | ALL [ PRIVILEGES ] }
ON FOREIGN DATA WRAPPER fdw_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]

{ USAGE | ALL [ PRIVILEGES ] }
ON FOREIGN SERVER server_name [, ...]
FROM { [ GROUP ] role_name | PUBLIC } [, ...]
[ CASCADE | RESTRICT ]
REVOKE

REVOKE [ GRANT OPTION FOR ]
  { EXECUTE | ALL [ PRIVILEGES ] }
  ON { FUNCTION function_name ( [ [ argmode ] [ arg_name ] arg_type [, ...] ] [, ...] | ALL FUNCTIONS IN SCHEMA schema_name [, ...] ) }
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { USAGE | ALL [ PRIVILEGES ] }
  ON LANGUAGE lang_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { { SELECT | UPDATE } [, ...] | ALL [ PRIVILEGES ] }
  ON LARGE OBJECT loid [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { { CREATE | USAGE } [, ...] | ALL [ PRIVILEGES ] }
  ON SCHEMA schema_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { CREATE | ALL [ PRIVILEGES ] }
  ON TABLESPACE tablespace_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ GRANT OPTION FOR ]
  { USAGE | ALL [ PRIVILEGES ] }
  ON TYPE type_name [, ...]
  FROM { [ GROUP ] role_name | PUBLIC } [, ...]
  [ CASCADE | RESTRICT ]

REVOKE [ ADMIN OPTION FOR ]
  role_name [, ...] FROM role_name [, ...]
  [ CASCADE | RESTRICT ]

Description

The REVOKE command revokes previously granted privileges from one or more roles. The key word PUBLIC refers to the implicitly defined group of all roles.

See the description of the GRANT command for the meaning of the privilege types.

Note that any particular role will have the sum of privileges granted directly to it, privileges granted to any role it is presently a member of, and privileges granted to PUBLIC. Thus, for example, revoking SELECT privilege from PUBLIC does not necessarily mean that all roles have lost SELECT privilege on the object; those who have it granted directly or via another role will still have it. Similarly, revoking SELECT from a user might not prevent that user from using SELECT if PUBLIC or another membership role still has SELECT rights.
If GRANT OPTION FOR is specified, only the grant option for the privilege is revoked, not the privilege itself. Otherwise, both the privilege and the grant option are revoked.

If a user holds a privilege with grant option and has granted it to other users then the privileges held by those other users are called dependent privileges. If the privilege or the grant option held by the first user is being revoked and dependent privileges exist, those dependent privileges are also revoked if CASCADE is specified; if it is not, the revoke action will fail. This recursive revocation only affects privileges that were granted through a chain of users that is traceable to the user that is the subject of this REVOKE command. Thus, the affected users might effectively keep the privilege if it was also granted through other users.

When revoking privileges on a table, the corresponding column privileges (if any) are automatically revoked on each column of the table, as well. On the other hand, if a role has been granted privileges on a table, then revoking the same privileges from individual columns will have no effect.

When revoking membership in a role, GRANT OPTION is instead called ADMIN OPTION, but the behavior is similar. Note also that this form of the command does not allow the noise word GROUP.

Notes

Use psql’s \dp command to display the privileges granted on existing tables and columns. See GRANT for information about the format. For non-table objects there are other \d commands that can display their privileges.

A user can only revoke privileges that were granted directly by that user. If, for example, user A has granted a privilege with grant option to user B, and user B has in turn granted it to user C, then user A cannot revoke the privilege directly from C. Instead, user A could revoke the grant option from user B and use the CASCADE option so that the privilege is in turn revoked from user C. For another example, if both A and B have granted the same privilege to C, A can revoke their own grant but not B’s grant, so C will still effectively have the privilege.

When a non-owner of an object attempts to REVOKE privileges on the object, the command will fail outright if the user has no privileges whatsoever on the object. As long as some privilege is available, the command will proceed, but it will revoke only those privileges for which the user has grant options. The REVOKE ALL PRIVILEGES forms will issue a warning message if no grant options are held, while the other forms will issue a warning if grant options for any of the privileges specifically named in the command are not held. (In principle these statements apply to the object owner as well, but since the owner is always treated as holding all grant options, the cases can never occur.)

If a superuser chooses to issue a GRANT or REVOKE command, the command is performed as though it were issued by the owner of the affected object. Since all privileges ultimately come from the object owner (possibly indirectly via chains of grant options), it is possible for a superuser to revoke all privileges, but this might require use of CASCADE as stated above.

REVOKE can also be done by a role that is not the owner of the affected object, but is a member of the role that owns the object, or is a member of a role that holds privileges WITH GRANT OPTION on the object. In this case the command is performed as though it were issued by the containing role that actually owns the object or holds the privileges WITH GRANT OPTION. For example, if table t1 is owned by role g1, of which role u1 is a member, then u1 can revoke privileges on t1 that are recorded as being granted by g1. This would include grants made by u1 as well as by other members of role g1.

If the role executing REVOKE holds privileges indirectly via more than one role membership path, it is unspecified which containing role will be used to perform the command. In such cases it is best practice to use SET ROLE to become the specific role you want to do the REVOKE as. Failure to do so might lead to revoking privileges other than the ones you intended, or not revoking anything at all.
Examples

Revoke insert privilege for the public on table films:
REVOKE INSERT ON films FROM PUBLIC;

Revoke all privileges from user manuel on view kinds:
REVOKE ALL PRIVILEGES ON kinds FROM manuel;

Note that this actually means “revoke all privileges that I granted”.
Revoke membership in role admins from user joe:
REVOKE admins FROM joe;

Compatibility

The compatibility notes of the GRANT command apply analogously to REVOKE. The keyword RESTRICT or CASCADE is required according to the standard, but PostgreSQL assumes RESTRICT by default.

See Also

GRANT
ROLLBACK

Name
ROLLBACK — abort the current transaction

Synopsis
ROLLBACK [ WORK | TRANSACTION ]

Description
ROLLBACK rolls back the current transaction and causes all the updates made by the transaction to be discarded.

Parameters
WORK
TRANSACTION
Optional key words. They have no effect.

Notes
Use COMMIT to successfully terminate a transaction.
Issuing ROLLBACK outside of a transaction block emits a warning and otherwise has no effect.

Examples
To abort all changes:
ROLLBACK;

Compatibility
The SQL standard only specifies the two forms ROLLBACK and ROLLBACK WORK. Otherwise, this command is fully conforming.

See Also
BEGIN, COMMIT, ROLLBACK TO SAVEPOINT
ROLLBACK PREPARED

Name
ROLLBACK PREPARED — cancel a transaction that was earlier prepared for two-phase commit

Synopsis
ROLLBACK PREPARED transaction_id

Description
ROLLBACK PREPARED rolls back a transaction that is in prepared state.

Parameters

transaction_id

The transaction identifier of the transaction that is to be rolled back.

Notes
To roll back a prepared transaction, you must be either the same user that executed the transaction originally, or a superuser. But you do not have to be in the same session that executed the transaction.

This command cannot be executed inside a transaction block. The prepared transaction is rolled back immediately.

All currently available prepared transactions are listed in the pg_prepared_xacts system view.

Examples
Roll back the transaction identified by the transaction identifier foobar:

ROLLBACK PREPARED 'foobar';

Compatibility
ROLLBACK PREPARED is a PostgreSQL extension. It is intended for use by external transaction management systems, some of which are covered by standards (such as X/Open XA), but the SQL side of those systems is not standardized.
See Also

PREPARE TRANSACTION, COMMIT PREPARED
ROLLBACK TO SAVEPOINT

Name

ROLLBACK TO SAVEPOINT — roll back to a savepoint

Synopsis

ROLLBACK [ WORK | TRANSACTION ] TO [ SAVEPOINT ] savepoint_name

Description

Roll back all commands that were executed after the savepoint was established. The savepoint remains valid and can be rolled back to again later, if needed.

ROLLBACK TO SAVEPOINT implicitly destroys all savepoints that were established after the named savepoint.

Parameters

savepoint_name

The savepoint to roll back to.

Notes

Use RELEASE SAVEPOINT to destroy a savepoint without discarding the effects of commands executed after it was established.

Specifying a savepoint name that has not been established is an error.

Cursors have somewhat non-transactional behavior with respect to savepoints. Any cursor that is opened inside a savepoint will be closed when the savepoint is rolled back. If a previously opened cursor is affected by a FETCH or MOVE command inside a savepoint that is later rolled back, the cursor remains at the position that FETCH left it pointing to (that is, the cursor motion caused by FETCH is not rolled back). Closing a cursor is not undone by rolling back, either. However, other side-effects caused by the cursor’s query (such as side-effects of volatile functions called by the query) are rolled back if they occur during a savepoint that is later rolled back. A cursor whose execution causes a transaction to abort is put in a cannot-execute state, so while the transaction can be restored using ROLLBACK TO SAVEPOINT, the cursor can no longer be used.

Examples

To undo the effects of the commands executed after my_savepoint was established:

ROLLBACK TO SAVEPOINT my_savepoint;
Cursor positions are not affected by savepoint rollback:

BEGIN;

DECLARE foo CURSOR FOR SELECT 1 UNION SELECT 2;

SAVEPOINT foo;

FETCH 1 FROM foo;
?column?
---------
  1

ROLLBACK TO SAVEPOINT foo;

FETCH 1 FROM foo;
?column?
---------
  2

COMMIT;

Compatibility

The SQL standard specifies that the key word SAVEPOINT is mandatory, but PostgreSQL and Oracle allow it to be omitted. SQL allows only WORK, not TRANSACTION, as a noise word after ROLLBACK. Also, SQL has an optional clause AND [ NO ] CHAIN which is not currently supported by PostgreSQL. Otherwise, this command conforms to the SQL standard.

See Also

BEGIN, COMMIT, RELEASE SAVEPOINT, ROLLBACK, SAVEPOINT
SAVEPOINT

Name
SAVEPOINT — define a new savepoint within the current transaction

Synopsis
SAVEPOINT savepoint_name

Description
SAVEPOINT establishes a new savepoint within the current transaction.
A savepoint is a special mark inside a transaction that allows all commands that are executed after it was established to be rolled back, restoring the transaction state to what it was at the time of the savepoint.

Parameters

savepoint_name

The name to give to the new savepoint.

Notes
Use ROLLBACK TO SAVEPOINT to rollback to a savepoint. Use RELEASE SAVEPOINT to destroy a savepoint, keeping the effects of commands executed after it was established.

Savepoints can only be established when inside a transaction block. There can be multiple savepoints defined within a transaction.

Examples
To establish a savepoint and later undo the effects of all commands executed after it was established:

BEGIN;
    INSERT INTO table1 VALUES (1);
    SAVEPOINT my_savepoint;
    INSERT INTO table1 VALUES (2);
    ROLLBACK TO SAVEPOINT my_savepoint;
    INSERT INTO table1 VALUES (3);
COMMIT;

The above transaction will insert the values 1 and 3, but not 2.

To establish and later destroy a savepoint:

BEGIN;
SAVEPOINT

```
INSERT INTO table1 VALUES (3);
SAVEPOINT my_savepoint;
INSERT INTO table1 VALUES (4);
RELEASE SAVEPOINT my_savepoint;
COMMIT;
```

The above transaction will insert both 3 and 4.

**Compatibility**

SQL requires a savepoint to be destroyed automatically when another savepoint with the same name is established. In PostgreSQL, the old savepoint is kept, though only the more recent one will be used when rolling back or releasing. (Releasing the newer savepoint with `RELEASE SAVEPOINT` will cause the older one to again become accessible.) Otherwise, `SAVEPOINT` is fully SQL conforming.

**See Also**

`BEGIN`, `COMMIT`, `RELEASE SAVEPOINT`, `ROLLBACK`, `ROLLBACK TO SAVEPOINT`
SECURITY LABEL

Name

SECURITY LABEL — define or change a security label applied to an object

Synopsis

SECURITY LABEL [ FOR provider ] ON
{
  TABLE object_name |
  COLUMN table_name.column_name |
  AGGREGATE aggregate_name ( aggregate_signature ) |
  DATABASE object_name |
  DOMAIN object_name |
  EVENT TRIGGER object_name |
  FOREIGN TABLE object_name
  FUNCTION function_name ( [ [ argmode ] argname argtype [, ... ] ] ) |
  LARGE OBJECT large_object_oid |
  MATERIALIZED VIEW object_name |
  [ PROCEDURAL ] LANGUAGE object_name |
  ROLE object_name |
  SCHEMA object_name |
  SEQUENCE object_name |
  TABLESPACE object_name |
  TYPE object_name |
  VIEW object_name
} IS 'label'

where aggregate_signature is:

* |
[ argmode ] | argname | argtype [, ... ] |
[ [ argmode ] argname argtype [, ... ] ] ORDER BY [ argmode ] | argname | argtype [, ... ] |

Description

SECURITY LABEL applies a security label to a database object. An arbitrary number of security labels, one per label provider, can be associated with a given database object. Label providers are loadable modules which register themselves by using the function register_label_provider.

Note: register_label_provider is not an SQL function; it can only be called from C code loaded into the backend.

The label provider determines whether a given label is valid and whether it is permissible to assign that label to a given object. The meaning of a given label is likewise at the discretion of the label provider. PostgreSQL places no restrictions on whether or how a label provider must interpret security labels; it merely provides a mechanism for storing them. In practice, this facility is intended to allow integration with label-based mandatory access control (MAC) systems such as SELinux. Such systems make all
access control decisions based on object labels, rather than traditional discretionary access control (DAC) concepts such as users and groups.

**Parameters**

- `object_name`
- `table_name.column_name`
- `aggregate_name`
- `function_name`

  The name of the object to be labeled. Names of tables, aggregates, domains, foreign tables, functions, sequences, types, and views can be schema-qualified.

- `provider`

  The name of the provider with which this label is to be associated. The named provider must be loaded and must consent to the proposed labeling operation. If exactly one provider is loaded, the provider name may be omitted for brevity.

- `argmode`

  The mode of a function or aggregate argument: `IN`, `OUT`, `INOUT`, or `VARIADIC`. If omitted, the default is `IN`. Note that `SECURITY LABEL` does not actually pay any attention to `OUT` arguments, since only the input arguments are needed to determine the function’s identity. So it is sufficient to list the `IN`, `INOUT`, and `VARIADIC` arguments.

- `argname`

  The name of a function or aggregate argument. Note that `SECURITY LABEL` does not actually pay any attention to argument names, since only the argument data types are needed to determine the function’s identity.

- `argtype`

  The data type of a function or aggregate argument.

- `large_object_oid`

  The OID of the large object.

**Examples**

The following example shows how the security label of a table might be changed.

```
SECURITY LABEL FOR selinux ON TABLE mytable IS 'system_u:object_r:sepgsql_table_t:s0';
```
Compatibility

There is no SECURITY LABEL command in the SQL standard.

See Also

sepgsql, src/test/modules/dummy_seclabel
SELECT

Name

SELECT, TABLE, WITH — retrieve rows from a table or view

Synopsis

[ WITH [ RECURSIVE ] with_query [, ... ] ]
SELECT [ ALL | DISTINCT [ ON ( expression [, ... ] ) ] ]

[ * | expression [ [ AS ] output_name ] [, ... ] ]
FROM from_item [, ... ]
WHERE condition
GROUP BY grouping_element [, ... ]
HAVING condition [, ... ]
WINDOW window_name AS ( window_definition ) [, ... ]
{ UNION | INTERSECT | EXCEPT } [ ALL | DISTINCT ] select
ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS { FIRST | LAST } ] [, ... ]
LIMIT { count | ALL }
OFFSET start [ ROW | ROWS ]
FETCH { FIRST | NEXT } [ count ] [ ROW | ROWS ] ONLY
FOR { UPDATE | NO KEY UPDATE | SHARE | KEY SHARE } [ OF table_name [, ... ] ] [ NOWAIT | SKIP LOCKED ]

where from_item can be one of:

[ ONLY ] table_name [ * ] [ [ AS ] alias [ ( column_alias [, ... ] ) ] ]
TABLESAMPLE sampling_method ( argument [, ... ] ) [ REPEATABLE ( seed ) ]
LATERAL ( select ) [ AS ] alias [ ( column_alias [, ... ] ) ]
with_query_name [ [ AS ] alias [ ( column_alias [, ... ] ) ] ]
LATERAL function_name ( [ argument [, ... ] ] )
[ WITH ORDINALITY ] [ [ AS ] alias [ ( column_alias [, ... ] ) ] ]
LATERAL function_name ( [ argument [, ... ] ] ) [ AS ] alias ( column_definition [, ... ] )
LATERAL function_name ( [ argument [, ... ] ] ) AS ( column_definition [, ... ] )
LATERAL ROWS FROM( function_name ( [ argument [, ... ] ] ) [ AS ] ( column_definition [ WITH ORDINALITY ] [ [ AS ] alias [ ( column_alias [, ... ] ) ] ] )
from_item [ NATURAL ] join_type from_item [ ON join_condition ] USING { join_column [, ... ]}

and grouping_element can be one of:

( )
expression
ROLLUP ( [ expression | ( expression [, ... ] ) ] [, ... ] )
CUBE ( [ expression | ( expression [, ... ] ) ] [, ... ] )
GROUPING SETS ( grouping_element [, ... ] )

and with_query is:

with_query_name [ ( column_name [, ... ] ) ] AS ( select | values | insert | update | delete)

TABLE [ ONLY ] table_name [ * ]

1711
**Description**

`SELECT` retrieves rows from zero or more tables. The general processing of `SELECT` is as follows:

1. All queries in the `WITH` list are computed. These effectively serve as temporary tables that can be referenced in the `FROM` list. A `WITH` query that is referenced more than once in `FROM` is computed only once. (See `WITH Clause` below.)

2. All elements in the `FROM` list are computed. (Each element in the `FROM` list is a real or virtual table.) If more than one element is specified in the `FROM` list, they are cross-joined together. (See `FROM Clause` below.)

3. If the `WHERE` clause is specified, all rows that do not satisfy the condition are eliminated from the output. (See `WHERE Clause` below.)

4. If the `GROUP BY` clause is specified, or if there are aggregate function calls, the output is combined into groups of rows that match on one or more values, and the results of aggregate functions are computed. If the `HAVING` clause is present, it eliminates groups that do not satisfy the given condition. (See `GROUP BY Clause` and `HAVING Clause` below.)

5. The actual output rows are computed using the `SELECT` output expressions for each selected row or row group. (See `SELECT List` below.)

6. `SELECT DISTINCT` eliminates duplicate rows from the result. `SELECT DISTINCT ON` eliminates rows that match on all the specified expressions. `SELECT ALL` (the default) will return all candidate rows, including duplicates. (See `DISTINCT Clause` below.)

7. Using the operators `UNION`, `INTERSECT`, and `EXCEPT`, the output of more than one `SELECT` statement can be combined to form a single result set. The `UNION` operator returns all rows that are in one or both of the result sets. The `INTERSECT` operator returns all rows that are strictly in both result sets. The `EXCEPT` operator returns the rows that are in the first result set but not in the second. In all three cases, duplicate rows are eliminated unless `ALL` is specified. The noise word `DISTINCT` can be added to explicitly specify eliminating duplicate rows. Notice that `DISTINCT` is the default behavior here, even though `ALL` is the default for `SELECT` itself. (See `UNION Clause`, `INTERSECT Clause`, and `EXCEPT Clause` below.)

8. If the `ORDER BY` clause is specified, the returned rows are sorted in the specified order. If `ORDER BY` is not given, the rows are returned in whatever order the system finds fastest to produce. (See `ORDER BY Clause` below.)

9. If the `LIMIT` (or `FETCH FIRST`) or `OFFSET` clause is specified, the `SELECT` statement only returns a subset of the result rows. (See `LIMIT Clause` below.)

10. If `FOR UPDATE`, `FOR NO KEY UPDATE`, `FOR SHARE` or `FOR KEY SHARE` is specified, the `SELECT` statement locks the selected rows against concurrent updates. (See `The Locking Clause` below.)

You must have `SELECT` privilege on each column used in a `SELECT` command. The use of `FOR NO KEY UPDATE`, `FOR UPDATE`, `FOR SHARE` or `FOR KEY SHARE` requires `UPDATE` privilege as well (for at least one column of each table so selected).
Parameters

WITH Clause

The WITH clause allows you to specify one or more subqueries that can be referenced by name in the primary query. The subqueries effectively act as temporary tables or views for the duration of the primary query. Each subquery can be a SELECT, TABLE, VALUES, INSERT, UPDATE or DELETE statement. When writing a data-modifying statement (INSERT, UPDATE or DELETE) in WITH, it is usual to include a RETURNING clause. It is the output of RETURNING, not the underlying table that the statement modifies, that forms the temporary table that is read by the primary query. If RETURNING is omitted, the statement is still executed, but it produces no output so it cannot be referenced as a table by the primary query.

A name (without schema qualification) must be specified for each WITH query. Optionally, a list of column names can be specified; if this is omitted, the column names are inferred from the subquery.

If RECURSIVE is specified, it allows a SELECT subquery to reference itself by name. Such a subquery must have the form

\[
\text{non_recursive_term} \text{ UNION [ ALL | DISTINCT ] recursive_term}
\]

where the recursive self-reference must appear on the right-hand side of the UNION. Only one recursive self-reference is permitted per query. Recursive data-modifying statements are not supported, but you can use the results of a recursive SELECT query in a data-modifying statement. See Section 7.8 for an example.

Another effect of RECURSIVE is that WITH queries need not be ordered: a query can reference another one that is later in the list. (However, circular references, or mutual recursion, are not implemented.) Without RECURSIVE, WITH queries can only reference sibling WITH queries that are earlier in the WITH list.

A key property of WITH queries is that they are evaluated only once per execution of the primary query, even if the primary query refers to them more than once. In particular, data-modifying statements are guaranteed to be executed once and only once, regardless of whether the primary query reads all or any of their output.

When there are multiple queries in the WITH clause, RECURSIVE should be written only once, immediately after WITH. It applies to all queries in the WITH clause, though it has no effect on queries that do not use recursion or forward references.

The primary query and the WITH queries are all (notionally) executed at the same time. This implies that the effects of a data-modifying statement in WITH cannot be seen from other parts of the query, other than by reading its RETURNING output. If two such data-modifying statements attempt to modify the same row, the results are unspecified.

See Section 7.8 for additional information.

FROM Clause

The FROM clause specifies one or more source tables for the SELECT. If multiple sources are specified, the result is the Cartesian product (cross join) of all the sources. But usually qualification conditions are added (via WHERE) to restrict the returned rows to a small subset of the Cartesian product.

The FROM clause can contain the following elements:
table_name

The name (optionally schema-qualified) of an existing table or view. If ONLY is specified before the table name, only that table is scanned. If ONLY is not specified, the table and all its descendant tables (if any) are scanned. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

alias

A substitute name for the FROM item containing the alias. An alias is used for brevity or to eliminate ambiguity for self-joins (where the same table is scanned multiple times). When an alias is provided, it completely hides the actual name of the table or function; for example given FROM foo AS f, the remainder of the SELECT must refer to this FROM item as f not foo. If an alias is written, a column alias list can also be written to provide substitute names for one or more columns of the table.

TABLESAMPLE sampling_method ( argument [, ...] ) [ REPEATABLE ( seed ) ]

A TABLESAMPLE clause after a table_name indicates that the specified sampling_method should be used to retrieve a subset of the rows in that table. This sampling precedes the application of any other filters such as WHERE clauses. The standard PostgreSQL distribution includes two sampling methods, BERNOULLI and SYSTEM, and other sampling methods can be installed in the database via extensions.

The BERNOULLI and SYSTEM sampling methods each accept a single argument which is the fraction of the table to sample, expressed as a percentage between 0 and 100. This argument can be any real-valued expression. (Other sampling methods might accept more or different arguments.) These two methods each return a randomly-chosen sample of the table that will contain approximately the specified percentage of the table's rows. The BERNOULLI method scans the whole table and selects or ignores individual rows independently with the specified probability. The SYSTEM method does block-level sampling with each block having the specified chance of being selected; all rows in each selected block are returned. The SYSTEM method is significantly faster than the BERNOULLI method when small sampling percentages are specified, but it may return a less-random sample of the table as a result of clustering effects.

The optional REPEATABLE clause specifies a seed number or expression to use for generating random numbers within the sampling method. The seed value can be any non-null floating-point value. Two queries that specify the same seed and argument values will select the same sample of the table, if the table has not been changed meanwhile. But different seed values will usually produce different samples. If REPEATABLE is not given then a new random sample is selected for each query, based upon a system-generated seed. Note that some add-on sampling methods do not accept REPEATABLE, and will always produce new samples on each use.

select

A sub-SELECT can appear in the FROM clause. This acts as though its output were created as a temporary table for the duration of this single SELECT command. Note that the sub-SELECT must be surrounded by parentheses, and an alias must be provided for it. A VALUES command can also be used here.

with_query_name

A WITH query is referenced by writing its name, just as though the query’s name were a table name. (In fact, the WITH query hides any real table of the same name for the purposes of the primary query. If necessary, you can refer to a real table of the same name by schema-qualifying the table’s name.) An alias can be provided in the same way as for a table.
**SELECT**

*function_name*

Function calls can appear in the **FROM** clause. (This is especially useful for functions that return result sets, but any function can be used.) This acts as though the function’s output were created as a temporary table for the duration of this single **SELECT** command. When the optional **WITH ORDINALITY** clause is added to the function call, a new column is appended after all the function’s output columns with numbering for each row.

An alias can be provided in the same way as for a table. If an alias is written, a column alias list can also be written to provide substitute names for one or more attributes of the function’s composite return type, including the column added by **ORDINALITY** if present.

Multiple function calls can be combined into a single **FROM**-clause item by surrounding them with **ROWS FROM( ... )**. The output of such an item is the concatenation of the first row from each function, then the second row from each function, etc. If some of the functions produce fewer rows than others, null values are substituted for the missing data, so that the total number of rows returned is always the same as for the function that produced the most rows.

If the function has been defined as returning the **record** data type, then an alias or the key word **AS** must be present, followed by a column definition list in the form `{ column_name data_type [, ... ] }`. The column definition list must match the actual number and types of columns returned by the function.

When using the **ROWS FROM( ... )** syntax, if one of the functions requires a column definition list, it’s preferred to put the column definition list after the function call inside **ROWS FROM( ... )**. A column definition list can be placed after the **ROWS FROM( ... )** construct only if there’s just a single function and no **WITH ORDINALITY** clause.

To use **ORDINALITY** together with a column definition list, you must use the **ROWS FROM( ... )** syntax and put the column definition list inside **ROWS FROM( ... )**.

**join_type**

One of

- [ INNER ] JOIN
- LEFT [ OUTER ] JOIN
- RIGHT [ OUTER ] JOIN
- FULL [ OUTER ] JOIN
- CROSS JOIN

For the **INNER** and **OUTER** join types, a join condition must be specified, namely exactly one of **NATURAL**, **ON join_condition**, or **USING (join_column [, ...])**. See below for the meaning. For **CROSS JOIN**, none of these clauses can appear.

A **JOIN** clause combines two **FROM** items, which for convenience we will refer to as “tables”, though in reality they can be any type of **FROM** item. Use parentheses if necessary to determine the order of nesting. In the absence of parentheses, **JOIN**s nest left-to-right. In any case **JOIN** binds more tightly than the commas separating **FROM**-list items.

**CROSS JOIN** and **INNER JOIN** produce a simple Cartesian product, the same result as you get from listing the two tables at the top level of **FROM**, but restricted by the join condition (if any). **CROSS JOIN** is equivalent to **INNER JOIN ON (TRUE)**, that is, no rows are removed by qualification. These join types are just a notational convenience, since they do nothing you couldn’t do with plain **FROM** and **WHERE**.

**LEFT OUTER JOIN** returns all rows in the qualified Cartesian product (i.e., all combined rows that pass its join condition), plus one copy of each row in the left-hand table for which there was
no right-hand row that passed the join condition. This left-hand row is extended to the full width
of the joined table by inserting null values for the right-hand columns. Note that only the JOIN
clause’s own condition is considered while deciding which rows have matches. Outer conditions
are applied afterwards.

Conversely, **RIGHT OUTER JOIN** returns all the joined rows, plus one row for each unmatched
right-hand row (extended with nulls on the left). This is just a notational convenience, since you
could convert it to a **LEFT OUTER JOIN** by switching the left and right tables.

**FULL OUTER JOIN** returns all the joined rows, plus one row for each unmatched left-hand row
(extended with nulls on the right), plus one row for each unmatched right-hand row (extended
with nulls on the left).

**ON** `join_condition`

`join_condition` is an expression resulting in a value of type `boolean` (similar to a **WHERE**
clause) that specifies which rows in a join are considered to match.

**USING** `{ join_column [, ...] }

A clause of the form **USING** `{ a, b, ... }` is shorthand for **ON** `left_table.a =
right_table.a AND left_table.b = right_table.b ...`. Also, **USING** implies that
only one of each pair of equivalent columns will be included in the join output, not both.

**NATURAL**

**NATURAL** is shorthand for a **USING** list that mentions all columns in the two tables that have
matching names. If there are no common column names, **NATURAL** is equivalent to **ON** `TRUE`.

**LATERAL**

The **LATERAL** key word can precede a sub-**SELECT** from item. This allows the sub-**SELECT** to
refer to columns of **FROM** items that appear before it in the **FROM** list. (Without **LATERAL**, each
sub-**SELECT** is evaluated independently and so cannot cross-reference any other **FROM** item.)

**LATERAL** can also precede a function-call **FROM** item, but in this case it is a noise word, because
the function expression can refer to earlier **FROM** items in any case.

A **LATERAL** item can appear at top level in the **FROM** list, or within a **JOIN** tree. In the latter case
it can also refer to any items that are on the left-hand side of a **JOIN** that it is on the right-hand
side of.

When a **FROM** item contains **LATERAL** cross-references, evaluation proceeds as follows: for each
row of the **FROM** item providing the cross-referenced column(s), or set of rows of multiple **FROM**
items providing the columns, the **LATERAL** item is evaluated using that row or row set’s values
of the columns. The resulting row(s) are joined as usual with the rows they were computed from.
This is repeated for each row or set of rows from the column source table(s).

The column source table(s) must be **INNER** or **LEFT** joined to the **LATERAL** item, else there would
not be a well-defined set of rows from which to compute each set of rows for the **LATERAL** item.
Thus, although a construct such as `X RIGHT JOIN LATERAL Y` is syntactically valid, it is not
actually allowed for `Y` to reference `X`.

**WHERE Clause**

The optional **WHERE** clause has the general form

**WHERE** `condition`
where condition is any expression that evaluates to a result of type boolean. Any row that does not satisfy this condition will be eliminated from the output. A row satisfies the condition if it returns true when the actual row values are substituted for any variable references.

**GROUP BY Clause**

The optional GROUP BY clause has the general form

```
GROUP BY grouping_element [, ...]
```

GROUP BY will condense into a single row all selected rows that share the same values for the grouped expressions. An expression used inside a grouping_element can be an input column name, or the name or ordinal number of an output column (SELECT list item), or an arbitrary expression formed from input-column values. In case of ambiguity, a GROUP BY name will be interpreted as an input-column name rather than an output column name.

If any of GROUPING SETS, ROLLUP or CUBE are present as grouping elements, then the GROUP BY clause as a whole defines some number of independent grouping sets. The effect of this is equivalent to constructing a UNION ALL between subqueries with the individual grouping sets as their GROUP BY clauses. For further details on the handling of grouping sets see Section 7.2.4.

Aggregate functions, if any are used, are computed across all rows making up each group, producing a separate value for each group. (If there are aggregate functions but no GROUP BY clause, the query is treated as having a single group comprising all the selected rows.) The set of rows fed to each aggregate function can be further filtered by attaching a FILTER clause to the aggregate function call; see Section 4.2.7 for more information. When a FILTER clause is present, only those rows matching it are included in the input to that aggregate function.

When GROUP BY is present, or any aggregate functions are present, it is not valid for the SELECT list expressions to refer to ungrouped columns except within aggregate functions or when the ungrouped column is functionally dependent on the grouped columns, since there would otherwise be more than one possible value to return for an ungrouped column. A functional dependency exists if the grouped columns (or a subset thereof) are the primary key of the table containing the ungrouped column.

Keep in mind that all aggregate functions are evaluated before evaluating any “scalar” expressions in the HAVING clause or SELECT list. This means that, for example, a CASE expression cannot be used to skip evaluation of an aggregate function; see Section 4.2.14.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified with GROUP BY.

**HAVING Clause**

The optional HAVING clause has the general form

```
HAVING condition
```

where condition is the same as specified for the WHERE clause.

HAVING eliminates group rows that do not satisfy the condition. HAVING is different from WHERE: WHERE filters individual rows before the application of GROUP BY, while HAVING filters group rows created by GROUP BY. Each column referenced in condition must unambiguously reference a grouping column, unless the reference appears within an aggregate function or the ungrouped column is functionally dependent on the grouping columns.
The presence of `HAVING` turns a query into a grouped query even if there is no `GROUP BY` clause. This is the same as what happens when the query contains aggregate functions but no `GROUP BY` clause. All the selected rows are considered to form a single group, and the `SELECT` list and `HAVING` clause can only reference table columns from within aggregate functions. Such a query will emit a single row if the `HAVING` condition is true, zero rows if it is not true.

Currently, `FOR NO KEY UPDATE`, `FOR UPDATE`, `FOR SHARE` and `FOR KEY SHARE` cannot be specified with `HAVING`.

**WINDOW Clause**

The optional `WINDOW` clause has the general form

```sql
WINDOW window_name AS ( window_definition ) [, ...]
```

where `window_name` is a name that can be referenced from `OVER` clauses or subsequent `window` definitions, and `window_definition` is

```sql
[ existing_window_name ]
[ PARTITION BY expression [, ...] ]
[ ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS { FIRST | LAST } ] [, ...] ]
[ frame_clause ]
```

If an `existing_window_name` is specified it must refer to an earlier entry in the `WINDOW` list; the new window copies its partitioning clause from that entry, as well as its ordering clause if any. In this case the new window cannot specify its own `PARTITION BY` clause, and it can specify `ORDER BY` only if the copied window does not have one. The new window always uses its own frame clause; the copied window must not specify a frame clause.

The elements of the `PARTITION BY` list are interpreted in much the same fashion as elements of a `GROUP BY Clause`, except that they are always simple expressions and never the name or number of an output column. Another difference is that these expressions can contain aggregate function calls, which are not allowed in a regular `GROUP BY` clause. They are allowed here because windowing occurs after grouping and aggregation.

Similarly, the elements of the `ORDER BY` list are interpreted in much the same fashion as elements of an `ORDER BY Clause`, except that the expressions are always taken as simple expressions and never the name or number of an output column.

The optional `frame_clause` defines the `window frame` for window functions that depend on the frame (not all do). The window frame is a set of related rows for each row of the query (called the `current row`). The `frame_clause` can be one of

```sql
{ RANGE | ROWS } frame_start
{ RANGE | ROWS } BETWEEN frame_start AND frame_end
```

where `frame_start` and `frame_end` can be one of

- `UNBOUNDED PRECEDING`
- `value PRECEDING`
- `CURRENT ROW`
- `value FOLLOWING`
- `UNBOUNDED FOLLOWING`
If frame_end is omitted it defaults to CURRENT ROW. Restrictions are that frame_start cannot be UNBOUNDED FOLLOWING, frame_end cannot be UNBOUNDED PRECEDING, and the frame_end choice cannot appear earlier in the above list than the frame_start choice — for example RANGE BETWEEN CURRENT ROW and value PRECEDING is not allowed.

The default framing option is RANGE UNBOUNDED PRECEDING, which is the same as RANGE BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW; it sets the frame to be all rows from the partition start up through the current row’s last peer (a row that ORDER BY considers equivalent to the current row, or all rows if there is no ORDER BY). In general, UNBOUNDED PRECEDING means that the frame starts with the first row of the partition, and similarly UNBOUNDED FOLLOWING means that the frame ends with the last row of the partition (regardless of RANGE or ROWS mode). In ROWS mode, CURRENT ROW means that the frame starts or ends with the current row; but in RANGE mode it means that the frame starts or ends with the current row’s first or last peer in the ORDER BY ordering.

The value PRECEDING and value FOLLOWING cases are currently only allowed in ROWS mode. They indicate that the frame starts or ends with the row that many rows before or after the current row. value must be an integer expression not containing any variables, aggregate functions, or window functions. The value must not be null or negative; but it can be zero, which selects the current row itself.

Beware that the ROWS options can produce unpredictable results if the ORDER BY ordering does not order the rows uniquely. The RANGE options are designed to ensure that rows that are peers in the ORDER BY ordering are treated alike; all peer rows will be in the same frame.

The purpose of a WINDOW clause is to specify the behavior of window functions appearing in the query’s SELECT List or ORDER BY Clause. These functions can reference the WINDOW clause entries by name in their OVER clauses. A WINDOW clause entry does not have to be referenced anywhere, however; if it is not used in the query it is simply ignored. It is possible to use window functions without any WINDOW clause at all, since a window function call can specify its window definition directly in its OVER clause. However, the WINDOW clause saves typing when the same window definition is needed for more than one window function.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified with WINDOW.

Window functions are described in detail in Section 3.5, Section 4.2.8, and Section 7.2.5.

**SELECT List**

The SELECT list (between the key words SELECT and FROM) specifies expressions that form the output rows of the SELECT statement. The expressions can (and usually do) refer to columns computed in the FROM clause.

Just as in a table, every output column of a SELECT has a name. In a simple SELECT this name is just used to label the column for display, but when the SELECT is a sub-query of a larger query, the name is seen by the larger query as the column name of the virtual table produced by the sub-query. To specify the name to use for an output column, write AS output_name after the column’s expression.

(You can omit AS, but only if the desired output name does not match any PostgreSQL keyword (see Appendix C). For protection against possible future keyword additions, it is recommended that you always either write AS or double-quote the output name.) If you do not specify a column name, a name is chosen automatically by PostgreSQL. If the column’s expression is a simple column reference then the chosen name is the same as that column’s name. In more complex cases a function or type name may be used, or the system may fall back on a generated name such as ?column?.

An output column’s name can be used to refer to the column’s value in ORDER BY and GROUP BY clauses, but not in the WHERE or HAVING clauses; there you must write out the expression instead.
Instead of an expression, * can be written in the output list as a shorthand for all the columns of the selected rows. Also, you can write `table_name.*` as a shorthand for the columns coming from just that table. In these cases it is not possible to specify new names with `AS`; the output column names will be the same as the table columns’ names.

According to the SQL standard, the expressions in the output list should be computed before applying `DISTINCT`, `ORDER BY`, or `LIMIT`. This is obviously necessary when using `DISTINCT`, since otherwise it’s not clear what values are being made distinct. However, in many cases it is convenient if output expressions are computed after `ORDER BY` and `LIMIT`; particularly if the output list contains any volatile or expensive functions. With that behavior, the order of function evaluations is more intuitive and there will not be evaluations corresponding to rows that never appear in the output. PostgreSQL will effectively evaluate output expressions after sorting and limiting, so long as those expressions are not referenced in `DISTINCT`, `ORDER BY` or `GROUP BY`. (As a counterexample, `SELECT f(x) FROM tab ORDER BY 1` clearly must evaluate `f(x)` before sorting.) Output expressions that contain set-returning functions are effectively evaluated after sorting and before limiting, so that `LIMIT` will act to cut off the output from a set-returning function.

**Note:** PostgreSQL versions before 9.6 did not provide any guarantees about the timing of evaluation of output expressions versus sorting and limiting; it depended on the form of the chosen query plan.

### DISTINCT Clause

If `SELECT DISTINCT` is specified, all duplicate rows are removed from the result set (one row is kept from each group of duplicates). `SELECT ALL` specifies the opposite: all rows are kept; that is the default.

`SELECT DISTINCT ON ( expression [, ...] )` keeps only the first row of each set of rows where the given expressions evaluate to equal. The `DISTINCT ON` expressions are interpreted using the same rules as for `ORDER BY` (see above). Note that the “first row” of each set is unpredictable unless `ORDER BY` is used to ensure that the desired row appears first. For example:

```
SELECT DISTINCT ON (location) location, time, report
  FROM weather_reports
  ORDER BY location, time DESC;
```

retrieves the most recent weather report for each location. But if we had not used `ORDER BY` to force descending order of time values for each location, we’d have gotten a report from an unpredictable time for each location.

The `DISTINCT ON` expression(s) must match the leftmost `ORDER BY` expression(s). The `ORDER BY` clause will normally contain additional expression(s) that determine the desired precedence of rows within each `DISTINCT ON` group.

Currently, `FOR NO KEY UPDATE`, `FOR UPDATE`, `FOR SHARE` and `FOR KEY SHARE` cannot be specified with `DISTINCT`.

### UNION Clause

The `UNION` clause has this general form:

```
select_statement UNION [ ALL | DISTINCT ] select_statement
```
**SELECT**

select_statement is any SELECT statement without an ORDER BY, LIMIT, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, or FOR KEY SHARE clause. (ORDER BY and LIMIT can be attached to a subexpression if it is enclosed in parentheses. Without parentheses, these clauses will be taken to apply to the result of the UNION, not to its right-hand input expression.)

The UNION operator computes the set union of the rows returned by the involved SELECT statements. A row is in the set union of two result sets if it appears in at least one of the result sets. The two SELECT statements that represent the direct operands of the UNION must produce the same number of columns, and corresponding columns must be of compatible data types.

The result of UNION does not contain any duplicate rows unless the ALL option is specified. ALL prevents elimination of duplicates. (Therefore, UNION ALL is usually significantly quicker than UNION; use ALL when you can.) DISTINCT can be written to explicitly specify the default behavior of eliminating duplicate rows.

Multiple UNION operators in the same SELECT statement are evaluated left to right, unless otherwise indicated by parentheses.

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified either for a UNION result or for any input of a UNION.

**INTERSECT** Clause

The INTERSECT clause has this general form:

```
select_statement INTERSECT [ ALL | DISTINCT ] select_statement
```

select_statement is any SELECT statement without an ORDER BY, LIMIT, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, or FOR KEY SHARE clause.

The INTERSECT operator computes the set intersection of the rows returned by the involved SELECT statements. A row is in the intersection of two result sets if it appears in both result sets.

The result of INTERSECT does not contain any duplicate rows unless the ALL option is specified. With ALL, a row that has m duplicates in the left table and n duplicates in the right table will appear min(m,n) times in the result set. DISTINCT can be written to explicitly specify the default behavior of eliminating duplicate rows.

Multiple INTERSECT operators in the same SELECT statement are evaluated left to right, unless parentheses dictate otherwise. INTERSECT binds more tightly than UNION. That is, A UNION B INTERSECT C will be read as A UNION (B INTERSECT C).

Currently, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE and FOR KEY SHARE cannot be specified either for an INTERSECT result or for any input of an INTERSECT.

**EXCEPT** Clause

The EXCEPT clause has this general form:

```
select_statement EXCEPT [ ALL | DISTINCT ] select_statement
```

select_statement is any SELECT statement without an ORDER BY, LIMIT, FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, or FOR KEY SHARE clause.

The EXCEPT operator computes the set of rows that are in the result of the left SELECT statement but not in the result of the right one.
The result of `EXCEPT` does not contain any duplicate rows unless the `ALL` option is specified. With `ALL`, a row that has \( m \) duplicates in the left table and \( n \) duplicates in the right table will appear \( \max(m-n,0) \) times in the result set. `DISTINCT` can be written to explicitly specify the default behavior of eliminating duplicate rows.

Multiple `EXCEPT` operators in the same `SELECT` statement are evaluated left to right, unless parentheses dictate otherwise. `EXCEPT` binds at the same level as `UNION`.

Currently, `FOR NO KEY UPDATE`, `FOR UPDATE`, `FOR SHARE` and `FOR KEY SHARE` cannot be specified either for an `EXCEPT` result or for any input of an `EXCEPT`.

**ORDER BY Clause**

The optional `ORDER BY` clause has this general form:

```
ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS { FIRST | LAST } ] [, ...]
```

The `ORDER BY` clause causes the result rows to be sorted according to the specified expression(s). If two rows are equal according to the leftmost expression, they are compared according to the next expression and so on. If they are equal according to all specified expressions, they are returned in an implementation-dependent order.

Each `expression` can be the name or ordinal number of an output column (`SELECT` list item), or it can be an arbitrary expression formed from input-column values.

The ordinal number refers to the ordinal (left-to-right) position of the output column. This feature makes it possible to define an ordering on the basis of a column that does not have a unique name. This is never absolutely necessary because it is always possible to assign a name to an output column using the `AS` clause.

It is also possible to use arbitrary expressions in the `ORDER BY` clause, including columns that do not appear in the `SELECT` output list. Thus the following statement is valid:

```
SELECT name FROM distributors ORDER BY code;
```

A limitation of this feature is that an `ORDER BY` clause applying to the result of a `UNION`, `INTERSECT`, or `EXCEPT` clause can only specify an output column name or number, not an expression.

If an `ORDER BY` expression is a simple name that matches both an output column name and an input column name, `ORDER BY` will interpret it as the output column name. This is the opposite of the choice that `GROUP BY` will make in the same situation. This inconsistency is made to be compatible with the SQL standard.

Optionally one can add the key word `ASC` (ascending) or `DESC` (descending) after any expression in the `ORDER BY` clause. If not specified, `ASC` is assumed by default. Alternatively, a specific ordering operator name can be specified in the `USING` clause. An ordering operator must be a less-than or greater-than member of some B-tree operator family. `ASC` is usually equivalent to `USING <` and `DESC` is usually equivalent to `USING >`. (But the creator of a user-defined data type can define exactly what the default sort ordering is, and it might correspond to operators with other names.)

If `NULLS LAST` is specified, null values sort after all non-null values; if `NULLS FIRST` is specified, null values sort before all non-null values. If neither is specified, the default behavior is `NULLS LAST` when `ASC` is specified or implied, and `NULLS FIRST` when `DESC` is specified (thus, the default is to act as though nulls are larger than non-nulls). When `USING` is specified, the default nulls ordering depends on whether the operator is a less-than or greater-than operator.

Note that ordering options apply only to the expression they follow; for example `ORDER BY x, y DESC` does not mean the same thing as `ORDER BY x DESC, y DESC`. 

1722
Character-string data is sorted according to the collation that applies to the column being sorted. That can be overridden at need by including a COLLATE clause in the expression, for example ORDER BY mycolumn COLLATE "en_US". For more information see Section 4.2.10 and Section 23.2.

**LIMIT Clause**

The **LIMIT** clause consists of two independent sub-clauses:

```
LIMIT { count | ALL }
OFFSET start
```

`count` specifies the maximum number of rows to return, while `start` specifies the number of rows to skip before starting to return rows. When both are specified, `start` rows are skipped before starting to count the `count` rows to be returned.

If the `count` expression evaluates to NULL, it is treated as **LIMIT ALL**, i.e., no limit. If `start` evaluates to NULL, it is treated the same as **OFFSET 0**.

SQL:2008 introduced a different syntax to achieve the same result, which PostgreSQL also supports. It is:

```
OFFSET start { ROW | ROWS }
FETCH { FIRST | NEXT } { count } { ROW | ROWS } ONLY
```

In this syntax, the `start` or `count` value is required by the standard to be a literal constant, a parameter, or a variable name; as a PostgreSQL extension, other expressions are allowed, but will generally need to be enclosed in parentheses to avoid ambiguity. If `count` is omitted in a **FETCH** clause, it defaults to 1. **ROW** and **ROWS** as well as **FIRST** and **NEXT** are noise words that don’t influence the effects of these clauses. According to the standard, the **OFFSET** clause must come before the **FETCH** clause if both are present; but PostgreSQL is laxer and allows either order.

When using **LIMIT**, it is a good idea to use an **ORDER BY** clause that constrains the result rows into a unique order. Otherwise you will get an unpredictable subset of the query’s rows — you might be asking for the tenth through twentieth rows, but tenth through twentieth in what ordering? You don’t know what ordering unless you specify **ORDER BY**.

The query planner takes **LIMIT** into account when generating a query plan, so you are very likely to get different plans (yielding different row orders) depending on what you use for **LIMIT** and **OFFSET**. Thus, using different **LIMIT**/**OFFSET** values to select different subsets of a query result will give *inconsistent results* unless you enforce a predictable result ordering with **ORDER BY**. This is not a bug; it is an inherent consequence of the fact that SQL does not promise to deliver the results of a query in any particular order unless **ORDER BY** is used to constrain the order.

It is even possible for repeated executions of the same **LIMIT** query to return different subsets of the rows of a table, if there is not an **ORDER BY** to enforce selection of a deterministic subset. Again, this is not a bug; determinism of the results is simply not guaranteed in such a case.

**The Locking Clause**

**FOR UPDATE**, **FOR NO KEY UPDATE**, **FOR SHARE** and **FOR KEY SHARE** are locking clauses; they affect how SELECT locks rows as they are obtained from the table.

The locking clause has the general form

```
FOR lock_strength [ OF table_name [, ... ] ] [ NOWAIT | SKIP LOCKED ]
```
where `lock_strength` can be one of

- `UPDATE`
- `NO KEY UPDATE`
- `SHARE`
- `KEY SHARE`

For more information on each row-level lock mode, refer to Section 13.3.2.

To prevent the operation from waiting for other transactions to commit, use either the `NOWAIT` or `SKIP LOCKED` option. With `NOWAIT`, the statement reports an error, rather than waiting, if a selected row cannot be locked immediately. With `SKIP LOCKED`, any selected rows that cannot be immediately locked are skipped. Skipping locked rows provides an inconsistent view of the data, so this is not suitable for general purpose work, but can be used to avoid lock contention with multiple consumers accessing a queue-like table. Note that `NOWAIT` and `SKIP LOCKED` apply only to the row-level lock(s) — the required `ROW SHARE` table-level lock is still taken in the ordinary way (see Chapter 13). You can use `LOCK` with the `NOWAIT` option first, if you need to acquire the table-level lock without waiting.

If specific tables are named in a locking clause, then only rows coming from those tables are locked; any other tables used in the `SELECT` are simply read as usual. A locking clause without a table list affects all tables used in the statement. If a locking clause is applied to a view or sub-query, it affects all tables used in the view or sub-query. However, these clauses do not apply to `WITH` queries referenced by the primary query. If you want row locking to occur within a `WITH` query, specify a locking clause within the `WITH` query.

Multiple locking clauses can be written if it is necessary to specify different locking behavior for different tables. If the same table is mentioned (or implicitly affected) by more than one locking clause, then it is processed as if it was only specified by the strongest one. Similarly, a table is processed as `NOWAIT` if that is specified in any of the clauses affecting it. Otherwise, it is processed as `SKIP LOCKED` if that is specified in any of the clauses affecting it.

The locking clauses cannot be used in contexts where returned rows cannot be clearly identified with individual table rows; for example they cannot be used with aggregation.

When a locking clause appears at the top level of a `SELECT` query, the rows that are locked are exactly those that are returned by the query; in the case of a join query, the rows locked are those that contribute to returned join rows. In addition, rows that satisfied the query conditions as of the query snapshot will be locked, although they will not be returned if they were updated after the snapshot and no longer satisfy the query conditions. If a `LIMIT` is used, locking stops once enough rows have been returned to satisfy the limit (but note that rows skipped over by `OFFSET` will get locked). Similarly, if a locking clause is used in a cursor’s query, only rows actually fetched or stepped past by the cursor will be locked.

When a locking clause appears in a sub-`SELECT`, the rows locked are those returned to the outer query by the sub-query. This might involve fewer rows than inspection of the sub-query alone would suggest, since conditions from the outer query might be used to optimize execution of the sub-query. For example,

```
SELECT * FROM (SELECT * FROM mytable FOR UPDATE) ss WHERE coll = 5;
```

will lock only rows having `coll = 5`, even though that condition is not textually within the sub-query.

Previous releases failed to preserve a lock which is upgraded by a later savepoint. For example, this code:
BEGIN;
SELECT * FROM mytable WHERE key = 1 FOR UPDATE;
SAVEPOINT s;
UPDATE mytable SET ... WHERE key = 1;
ROLLBACK TO s;

would fail to preserve the FOR UPDATE lock after the ROLLBACK TO. This has been fixed in release 9.3.

Caution

It is possible for a SELECT command running at the READ COMMITTED transaction isolation level and using ORDER BY and a locking clause to return rows out of order. This is because ORDER BY is applied first. The command sorts the result, but might then block trying to obtain a lock on one or more of the rows. Once the SELECT unblocks, some of the ordering column values might have been modified, leading to those rows appearing to be out of order (though they are in order in terms of the original column values). This can be worked around at need by placing the FOR UPDATE/SHARE clause in a sub-query, for example

```sql
SELECT * FROM (SELECT * FROM mytable FOR UPDATE) ss ORDER BY column1;
```

Note that this will result in locking all rows of mytable, whereas FOR UPDATE at the top level would lock only the actually returned rows. This can make for a significant performance difference, particularly if the ORDER BY is combined with LIMIT or other restrictions. So this technique is recommended only if concurrent updates of the ordering columns are expected and a strictly sorted result is required.

At the REPEATABLE READ or SERIALIZABLE transaction isolation level this would cause a serialization failure (with a SQLSTATE of ‘40001’), so there is no possibility of receiving rows out of order under these isolation levels.

TABLE Command

The command

```
TABLE name
```

is equivalent to

```
SELECT * FROM name
```

It can be used as a top-level command or as a space-saving syntax variant in parts of complex queries. Only the WITH, UNION, INTERSECT, EXCEPT, ORDER BY, LIMIT, OFFSET, FETCH and FOR locking clauses can be used with TABLE; the WHERE clause and any form of aggregation cannot be used.

Examples

To join the table films with the table distributors:

```sql
SELECT f.title, f.did, d.name, f.date_prod, f.kind
  FROM distributors d, films f
  WHERE f.did = d.did
```
To sum the column `len` of all films and group the results by `kind`:

```
SELECT kind, sum(len) AS total FROM films GROUP BY kind;
```

<table>
<thead>
<tr>
<th>kind</th>
<th>total</th>
</tr>
</thead>
<tbody>
<tr>
<td>Action</td>
<td>07:34</td>
</tr>
<tr>
<td>Comedy</td>
<td>02:58</td>
</tr>
<tr>
<td>Drama</td>
<td>14:28</td>
</tr>
<tr>
<td>Musical</td>
<td>06:42</td>
</tr>
<tr>
<td>Romantic</td>
<td>04:38</td>
</tr>
</tbody>
</table>

To sum the column `len` of all films, group the results by `kind` and show those group totals that are less than 5 hours:

```
SELECT kind, sum(len) AS total
FROM films
GROUP BY kind
HAVING sum(len) < interval '5 hours';
```

<table>
<thead>
<tr>
<th>kind</th>
<th>total</th>
</tr>
</thead>
<tbody>
<tr>
<td>Comedy</td>
<td>02:58</td>
</tr>
<tr>
<td>Romantic</td>
<td>04:38</td>
</tr>
</tbody>
</table>

The following two examples are identical ways of sorting the individual results according to the contents of the second column (`name`):

```
SELECT * FROM distributors ORDER BY name;
SELECT * FROM distributors ORDER BY 2;
```

<table>
<thead>
<tr>
<th>did</th>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>109</td>
<td>20th Century Fox</td>
</tr>
<tr>
<td>110</td>
<td>Bavaria Atelier</td>
</tr>
<tr>
<td>101</td>
<td>British Lion</td>
</tr>
<tr>
<td>107</td>
<td>Columbia</td>
</tr>
<tr>
<td>102</td>
<td>Jean Luc Godard</td>
</tr>
<tr>
<td>113</td>
<td>Luso films</td>
</tr>
<tr>
<td>104</td>
<td>Mosfilm</td>
</tr>
<tr>
<td>103</td>
<td>Paramount</td>
</tr>
<tr>
<td>106</td>
<td>Toho</td>
</tr>
<tr>
<td>105</td>
<td>United Artists</td>
</tr>
<tr>
<td>111</td>
<td>Walt Disney</td>
</tr>
<tr>
<td>112</td>
<td>Warner Bros.</td>
</tr>
<tr>
<td>108</td>
<td>Westward</td>
</tr>
</tbody>
</table>
The next example shows how to obtain the union of the tables `distributors` and `actors`, restricting the results to those that begin with the letter W in each table. Only distinct rows are wanted, so the key word `ALL` is omitted.

<table>
<thead>
<tr>
<th>did</th>
<th>name</th>
<th>id</th>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>108</td>
<td>Westward</td>
<td>1</td>
<td>Woody Allen</td>
</tr>
<tr>
<td>111</td>
<td>Walt Disney</td>
<td>2</td>
<td>Warren Beatty</td>
</tr>
<tr>
<td>112</td>
<td>Warner Bros.</td>
<td>3</td>
<td>Walter Matthau</td>
</tr>
</tbody>
</table>

```
SELECT distributors.name
FROM distributors
WHERE distributors.name LIKE 'W%'
UNION
SELECT actors.name
FROM actors
WHERE actors.name LIKE 'W%';
```

This example shows how to use a function in the `FROM` clause, both with and without a column definition list:

```
CREATE FUNCTION distributors(int) RETURNS SETOF distributors AS $$
SELECT * FROM distributors WHERE did = $1;
$$ LANGUAGE SQL;

SELECT * FROM distributors(111);
```

```
CREATE FUNCTION distributors_2(int) RETURNS SETOF record AS $$
SELECT * FROM distributors WHERE did = $1;
$$ LANGUAGE SQL;

SELECT * FROM distributors_2(111) AS (f1 int, f2 text);
```

Here is an example of a function with an ordinality column added:

```
SELECT * FROM unnest(ARRAY['a','b','c','d','e','f']) WITH ORDINALITY;
```
This example shows how to use a simple WITH clause:

WITH t AS (  
    SELECT random() as x FROM generate_series(1, 3)  
)  
SELECT * FROM t  
UNION ALL  
SELECT * FROM t

Notice that the WITH query was evaluated only once, so that we got two sets of the same three random values.

This example uses WITH RECURSIVE to find all subordinates (direct or indirect) of the employee Mary, and their level of indirectness, from a table that shows only direct subordinates:

WITH RECURSIVE employee_recursive(distance, employee_name, manager_name) AS (  
    SELECT 1, employee_name, manager_name  
    FROM employee  
    WHERE manager_name = 'Mary'  
    UNION ALL  
    SELECT er.distance + 1, e.employee_name, e.manager_name  
    FROM employee_recursive er, employee e  
    WHERE er.employee_name = e.manager_name  
)  
SELECT distance, employee_name FROM employee_recursive;

Notice the typical form of recursive queries: an initial condition, followed by UNION, followed by the recursive part of the query. Be sure that the recursive part of the query will eventually return no tuples, or else the query will loop indefinitely. (See Section 7.8 for more examples.)

This example uses LATERAL to apply a set-returning function get_product_names() for each row of the manufacturers table:

SELECT m.name AS mname, pname  
FROM manufacturers m, LATERAL get_product_names(m.id) pname;

Manufacturers not currently having any products would not appear in the result, since it is an inner join. If we wished to include the names of such manufacturers in the result, we could do:
**SELECT**

```sql
SELECT m.name AS mname, pname 
FROM manufacturers m LEFT JOIN LATERAL get_product_names(m.id) pname ON true;
```

**Compatibility**

Of course, the `SELECT` statement is compatible with the SQL standard. But there are some extensions and some missing features.

**Omitted FROM Clauses**

PostgreSQL allows one to omit the `FROM` clause. It has a straightforward use to compute the results of simple expressions:

```sql
SELECT 2+2;
```

<table>
<thead>
<tr>
<th>?column?</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
</tr>
</tbody>
</table>

Some other SQL databases cannot do this except by introducing a dummy one-row table from which to do the `SELECT`.

Note that if a `FROM` clause is not specified, the query cannot reference any database tables. For example, the following query is invalid:

```sql
SELECT distributors.* WHERE distributors.name = 'Westward';
```

PostgreSQL releases prior to 8.1 would accept queries of this form, and add an implicit entry to the query’s `FROM` clause for each table referenced by the query. This is no longer allowed.

**Empty SELECT Lists**

The list of output expressions after `SELECT` can be empty, producing a zero-column result table. This is not valid syntax according to the SQL standard. PostgreSQL allows it to be consistent with allowing zero-column tables. However, an empty list is not allowed when `DISTINCT` is used.

**Omitting the AS Key Word**

In the SQL standard, the optional key word `AS` can be omitted before an output column name whenever the new column name is a valid column name (that is, not the same as any reserved keyword). PostgreSQL is slightly more restrictive: `AS` is required if the new column name matches any keyword at all, reserved or not. Recommended practice is to use `AS` or double-quote output column names, to prevent any possible conflict against future keyword additions.

In `FROM` items, both the standard and PostgreSQL allow `AS` to be omitted before an alias that is an unreserved keyword. But this is impractical for output column names, because of syntactic ambiguities.
**ONLY and Inheritance**

The SQL standard requires parentheses around the table name when writing `ONLY`, for example

```
SELECT * FROM ONLY (tab1), ONLY (tab2) WHERE ... .
```

PostgreSQL considers these parentheses to be optional.

PostgreSQL allows a trailing `*` to be written to explicitly specify the non-ONLY behavior of including child tables. The standard does not allow this.

(These points apply equally to all SQL commands supporting the `ONLY` option.)

**TABLESAMPLE Clause Restrictions**

The `TABLESAMPLE` clause is currently accepted only on regular tables and materialized views. According to the SQL standard it should be possible to apply it to any `FROM` item.

**Function Calls in FROM**

PostgreSQL allows a function call to be written directly as a member of the `FROM` list. In the SQL standard it would be necessary to wrap such a function call in a sub-SELECT; that is, the syntax `FROM func(...) alias` is approximately equivalent to `FROM LATERAL (SELECT func(...)) alias`. Note that `LATERAL` is considered to be implicit; this is because the standard requires `LATERAL` semantics for an `UNNEST()` item in `FROM`. PostgreSQL treats `UNNEST()` the same as other set-returning functions.

**Namespace Available to GROUP BY and ORDER BY**

In the SQL-92 standard, an `ORDER BY` clause can only use output column names or numbers, while a `GROUP BY` clause can only use expressions based on input column names. PostgreSQL extends each of these clauses to allow the other choice as well (but it uses the standard’s interpretation if there is ambiguity). PostgreSQL also allows both clauses to specify arbitrary expressions. Note that names appearing in an expression will always be taken as input-column names, not as output-column names.

SQL:1999 and later use a slightly different definition which is not entirely upward compatible with SQL-92. In most cases, however, PostgreSQL will interpret an `ORDER BY` or `GROUP BY` expression the same way SQL:1999 does.

**Functional Dependencies**

PostgreSQL recognizes functional dependency (allowing columns to be omitted from `GROUP BY`) only when a table’s primary key is included in the `GROUP BY` list. The SQL standard specifies additional conditions that should be recognized.

**WINDOW Clause Restrictions**

The SQL standard provides additional options for the window `frame_clause`. PostgreSQL currently supports only the options listed above.
**SELECT**

**LIMIT and OFFSET**

The clauses `LIMIT` and `OFFSET` are PostgreSQL-specific syntax, also used by MySQL. The SQL:2008 standard has introduced the clauses `OFFSET ... FETCH {FIRST|NEXT} ...` for the same functionality, as shown above in *LIMIT Clause*. This syntax is also used by IBM DB2.

(Applications written for Oracle frequently use a workaround involving the automatically generated `rownum` column, which is not available in PostgreSQL, to implement the effects of these clauses.)

**FOR NO KEY UPDATE, FOR UPDATE, FOR SHARE, FOR KEY SHARE**

Although `FOR UPDATE` appears in the SQL standard, the standard allows it only as an option of `DECLARE CURSOR`. PostgreSQL allows it in any `SELECT` query as well as in sub-`SELECTs`, but this is an extension. The `FOR NO KEY UPDATE, FOR SHARE` and `FOR KEY SHARE` variants, as well as the `NOWAIT` and `SKIP LOCKED` options, do not appear in the standard.

**Data-Modifying Statements in WITH**

PostgreSQL allows `INSERT`, `UPDATE`, and `DELETE` to be used as `WITH` queries. This is not found in the SQL standard.

**Nonstandard Clauses**

`DISTINCT ON ( ... )` is an extension of the SQL standard.

`ROWS FROM( ... )` is an extension of the SQL standard.
**SELECT INTO**

**Name**

SELECT INTO — define a new table from the results of a query

**Synopsis**

```sql
[ WITH [ RECURSIVE ] with_query [, ... ] ]
SELECT [ ALL | DISTINCT [ ON ( expression [, ... ] ) ] ]
  * | expression [ AS output_name ] [, ...]
  INTO [ TEMPORARY | TEMP | UNLOGGED ] [ TABLE ] new_table
  [ FROM from_item [, ...] ]
  [ WHERE condition ]
  [ GROUP BY expression [, ...] ]
  [ HAVING condition [, ...] ]
  [ WINDOW window_name AS ( window_definition ) [, ...] ]
  [ { UNION | INTERSECT | EXCEPT } [ ALL | DISTINCT ] select ]
  [ ORDER BY expression [ ASC | DESC | USING operator ] [ NULLS { FIRST | LAST } ] [, ...]
  LIMIT { count | ALL } ]
  [ OFFSET start [ ROW | ROWS ] ]
  [ FETCH { FIRST | NEXT } { count | ROW | ROWS } ONLY ]
  [ FOR { UPDATE | SHARE } { OF table_name [, ...] | NOWAIT } [, ...] ]
```

**Description**

**SELECT INTO** creates a new table and fills it with data computed by a query. The data is not returned to the client, as it is with a normal **SELECT**. The new table’s columns have the names and data types associated with the output columns of the **SELECT**.

**Parameters**

**TEMPORARY** or **TEMP**

If specified, the table is created as a temporary table. Refer to **CREATE TABLE** for details.

**UNLOGGED**

If specified, the table is created as an unlogged table. Refer to **CREATE TABLE** for details.

**new_table**

The name (optionally schema-qualified) of the table to be created.

All other parameters are described in detail under **SELECT**.

**Notes**

**CREATE TABLE AS** is functionally similar to **SELECT INTO**. **CREATE TABLE AS** is the recommended syntax, since this form of **SELECT INTO** is not available in **ECPG** or **PL/pgSQL**, because
they interpret the INTO clause differently. Furthermore, CREATE TABLE AS offers a superset of the functionality provided by SELECT INTO.

To add OIDs to the table created by SELECT INTO, enable the default_with_oids configuration variable. Alternatively, CREATE TABLE AS can be used with the WITH OIDS clause.

Examples

Create a new table films_recent consisting of only recent entries from the table films:

```
SELECT * INTO films_recent FROM films WHERE date_prod > '2002-01-01';
```

Compatibility

The SQL standard uses SELECT INTO to represent selecting values into scalar variables of a host program, rather than creating a new table. This indeed is the usage found in ECPG (see Chapter 34) and PL/pgSQL (see Chapter 41). The PostgreSQL usage of SELECT INTO to represent table creation is historical. It is best to use CREATE TABLE AS for this purpose in new code.

See Also

CREATE TABLE AS
SET

Name

SET — change a run-time parameter

Synopsis

SET [ SESSION | LOCAL ] configuration_parameter { TO | = } { value | 'value' | DEFAULT }

SET [ SESSION | LOCAL ] TIME ZONE { timezone | LOCAL | DEFAULT }

Description

The SET command changes run-time configuration parameters. Many of the run-time parameters listed in Chapter 19 can be changed on-the-fly with SET. (But some require superuser privileges to change, and others cannot be changed after server or session start.) SET only affects the value used by the current session.

If SET (or equivalently SET SESSION) is issued within a transaction that is later aborted, the effects of the SET command disappear when the transaction is rolled back. Once the surrounding transaction is committed, the effects will persist until the end of the session, unless overridden by another SET.

The effects of SET LOCAL last only till the end of the current transaction, whether committed or not. A special case is SET followed by SET LOCAL within a single transaction: the SET LOCAL value will be seen until the end of the transaction, but afterwards (if the transaction is committed) the SET value will take effect.

The effects of SET or SET LOCAL are also canceled by rolling back to a savepoint that is earlier than the command.

If SET LOCAL is used within a function that has a SET option for the same variable (see CREATE FUNCTION), the effects of the SET LOCAL command disappear at function exit; that is, the value in effect when the function was called is restored anyway. This allows SET LOCAL to be used for dynamic or repeated changes of a parameter within a function, while still having the convenience of using the SET option to save and restore the caller’s value. However, a regular SET command overrides any surrounding function’s SET option; its effects will persist unless rolled back.

Note: In PostgreSQL versions 8.0 through 8.2, the effects of a SET LOCAL would be canceled by releasing an earlier savepoint, or by successful exit from a PL/pgSQL exception block. This behavior has been changed because it was deemed unintuitive.

Parameters

SESSION

Specifies that the command takes effect for the current session. (This is the default if neither SESSION nor LOCAL appears.)
SET LOCAL

Specifies that the command takes effect for only the current transaction. After COMMIT or ROLLBACK, the session-level setting takes effect again. Issuing this outside of a transaction block emits a warning and otherwise has no effect.

coloration parameter

Name of a settable run-time parameter. Available parameters are documented in Chapter 19 and below.

value

New value of parameter. Values can be specified as string constants, identifiers, numbers, or comma-separated lists of these, as appropriate for the particular parameter. DEFAULT can be written to specify resetting the parameter to its default value (that is, whatever value it would have had if no SET had been executed in the current session).

Besides the configuration parameters documented in Chapter 19, there are a few that can only be adjusted using the SET command or that have a special syntax:

SCHEMA

SET SCHEMA 'value' is an alias for SET search_path TO value. Only one schema can be specified using this syntax.

NAMES

SET NAMES value is an alias for SET client_encoding TO value.

SEED

Sets the internal seed for the random number generator (the function random). Allowed values are floating-point numbers between -1 and 1, which are then multiplied by $2^{31}-1$.

The seed can also be set by invoking the function setseed:

SELECT setseed(value);

TIME ZONE

SET TIME ZONE value is an alias for SET timezone TO value. The syntax SET TIME ZONE allows special syntax for the time zone specification. Here are examples of valid values:

'PST8PDT'

The time zone for Berkeley, California.

'Europe/Rome'

The time zone for Italy.

-7

The time zone 7 hours west from UTC (equivalent to PDT). Positive values are east from UTC.

INTERVAL '-08:00' HOUR TO MINUTE

The time zone 8 hours west from UTC (equivalent to PST).

LOCAL

Set the time zone to your local time zone (that is, the server's default value of timezone).
Timezone settings given as numbers or intervals are internally translated to POSIX timezone syntax. For example, after `SET TIME ZONE -7`, `SHOW TIME ZONE` would report `<-07>+07`. See Section 8.5.3 for more information about time zones.

**Notes**

The function `set_config` provides equivalent functionality; see Section 9.26. Also, it is possible to `UPDATE` the `pg_settings` system view to perform the equivalent of `SET`.

**Examples**

Set the schema search path:

```
SET search_path TO my_schema, public;
```

Set the style of date to traditional POSTGRES with “day before month” input convention:

```
SET datestyle TO postgres, dmy;
```

Set the time zone for Berkeley, California:

```
SET TIME ZONE 'PST8PDT';
```

Set the time zone for Italy:

```
SET TIME ZONE 'Europe/Rome';
```

**Compatibility**

`SET TIME ZONE` extends syntax defined in the SQL standard. The standard allows only numeric time zone offsets while PostgreSQL allows more flexible time-zone specifications. All other `SET` features are PostgreSQL extensions.

**See Also**

`RESET`, `SHOW`
SET CONSTRAINTS

Name
SET CONSTRAINTS — set constraint check timing for the current transaction

Synopsis
SET CONSTRAINTS { ALL | name [, ...] } { DEFERRED | IMMEDIATE }

Description
SET CONSTRAINTS sets the behavior of constraint checking within the current transaction. IMMEDIATE constraints are checked at the end of each statement. DEFERRED constraints are not checked until transaction commit. Each constraint has its own IMMEDIATE or DEFERRED mode.

Upon creation, a constraint is given one of three characteristics: DEFERRABLE INITIALLY DEFERRED, DEFERRABLE INITIALLY IMMEDIATE, or NOT DEFERRABLE. The third class is always IMMEDIATE and is not affected by the SET CONSTRAINTS command. The first two classes start every transaction in the indicated mode, but their behavior can be changed within a transaction by SET CONSTRAINTS.

SET CONSTRAINTS with a list of constraint names changes the mode of just those constraints (which must all be deferrable). Each constraint name can be schema-qualified. The current schema search path is used to find the first matching name if no schema name is specified. SET CONSTRAINTS ALL changes the mode of all deferrable constraints.

When SET CONSTRAINTS changes the mode of a constraint from DEFERRED to IMMEDIATE, the new mode takes effect retroactively: any outstanding data modifications that would have been checked at the end of the transaction are instead checked during the execution of the SET CONSTRAINTS command. If any such constraint is violated, the SET CONSTRAINTS fails (and does not change the constraint mode). Thus, SET CONSTRAINTS can be used to force checking of constraints to occur at a specific point in a transaction.

Currently, only UNIQUE, PRIMARY KEY, REFERENCES (foreign key), and EXCLUDE constraints are affected by this setting. NOT NULL and CHECK constraints are always checked immediately when a row is inserted or modified (not at the end of the statement). Uniqueness and exclusion constraints that have not been declared DEFERRABLE are also checked immediately.

The firing of triggers that are declared as “constraint triggers” is also controlled by this setting — they fire at the same time that the associated constraint should be checked.

Notes
Because PostgreSQL does not require constraint names to be unique within a schema (but only per-table), it is possible that there is more than one match for a specified constraint name. In this case SET CONSTRAINTS will act on all matches. For a non-schema-qualified name, once a match or matches have been found in some schema in the search path, schemas appearing later in the path are not searched.

This command only alters the behavior of constraints within the current transaction. Issuing this outside of a transaction block emits a warning and otherwise has no effect.
Compatibility

This command complies with the behavior defined in the SQL standard, except for the limitation that, in PostgreSQL, it does not apply to NOT NULL and CHECK constraints. Also, PostgreSQL checks non-deferrable uniqueness constraints immediately, not at end of statement as the standard would suggest.
**SET ROLE**

**Name**

SET ROLE — set the current user identifier of the current session

**Synopsis**

```sql
SET [ SESSION | LOCAL ] ROLE role_name
SET [ SESSION | LOCAL ] ROLE NONE
RESET ROLE
```

**Description**

This command sets the current user identifier of the current SQL session to be `role_name`. The role name can be written as either an identifier or a string literal. After `SET ROLE`, permissions checking for SQL commands is carried out as though the named role were the one that had logged in originally.

The specified `role_name` must be a role that the current session user is a member of. (If the session user is a superuser, any role can be selected.)

The `SESSION` and `LOCAL` modifiers act the same as for the regular `SET` command.

The `NONE` and `RESET` forms reset the current user identifier to be the current session user identifier. These forms can be executed by any user.

**Notes**

Using this command, it is possible to either add privileges or restrict one’s privileges. If the session user role has the `INHERIT` attribute, then it automatically has all the privileges of every role that it could SET ROLE to; in this case SET ROLE effectively drops all the privileges assigned directly to the session user and to the other roles it is a member of, leaving only the privileges available to the named role. On the other hand, if the session user role has the `NOINHERIT` attribute, SET ROLE drops the privileges assigned directly to the session user and instead acquires the privileges available to the named role.

In particular, when a superuser chooses to SET ROLE to a non-superuser role, they lose their superuser privileges.

`SET ROLE` has effects comparable to `SET SESSION AUTHORIZATION`, but the privilege checks involved are quite different. Also, `SET SESSION AUTHORIZATION` determines which roles are allowable for later `SET ROLE` commands, whereas changing roles with `SET ROLE` does not change the set of roles allowed to a later `SET ROLE`.

`SET ROLE` does not process session variables as specified by the role’s `ALTER ROLE` settings; this only happens during login.

`SET ROLE` cannot be used within a `SECURITY DEFINER` function.
Examples

```sql
SELECT SESSION_USER, CURRENT_USER;

session_user  |  current_user
--------------+--------------
peter         |  peter

SET ROLE 'paul';

SELECT SESSION_USER, CURRENT_USER;

session_user  |  current_user
--------------+--------------
peter         |  paul
```

Compatibility

PostgreSQL allows identifier syntax ("rolename"), while the SQL standard requires the role name to be written as a string literal. SQL does not allow this command during a transaction; PostgreSQL does not make this restriction because there is no reason to. The `SESSION` and `LOCAL` modifiers are a PostgreSQL extension, as is the `RESET` syntax.

See Also

`SET SESSION AUTHORIZATION`
SET SESSION AUTHORIZATION

Name

SET SESSION AUTHORIZATION — set the session user identifier and the current user identifier of the current session

Synopsis

SET [ SESSION | LOCAL ] SESSION AUTHORIZATION  user_name
SET [ SESSION | LOCAL ] SESSION AUTHORIZATION  DEFAULT
RESET SESSION AUTHORIZATION

Description

This command sets the session user identifier and the current user identifier of the current SQL session to be user_name. The user name can be written as either an identifier or a string literal. Using this command, it is possible, for example, to temporarily become an unprivileged user and later switch back to being a superuser.

The session user identifier is initially set to be the (possibly authenticated) user name provided by the client. The current user identifier is normally equal to the session user identifier, but might change temporarily in the context of SECURITY DEFINER functions and similar mechanisms; it can also be changed by SET ROLE. The current user identifier is relevant for permission checking.

The session user identifier can be changed only if the initial session user (the authenticated user) had the superuser privilege. Otherwise, the command is accepted only if it specifies the authenticated user name.

The SESSION and LOCAL modifiers act the same as for the regular SET command.
The DEFAULT and RESET forms reset the session and current user identifiers to be the originally authenticated user name. These forms can be executed by any user.

Notes

SET SESSION AUTHORIZATION cannot be used within a SECURITY DEFINER function.

Examples

SELECT SESSION_USER, CURRENT_USER;

<table>
<thead>
<tr>
<th>session_user</th>
<th>current_user</th>
</tr>
</thead>
<tbody>
<tr>
<td>peter</td>
<td>peter</td>
</tr>
</tbody>
</table>

SET SESSION AUTHORIZATION  'paul';

SELECT SESSION_USER, CURRENT_USER;
Compatibility

The SQL standard allows some other expressions to appear in place of the literal user_name, but these options are not important in practice. PostgreSQL allows identifier syntax ("username"), which SQL does not. SQL does not allow this command during a transaction; PostgreSQL does not make this restriction because there is no reason to. The SESSION and LOCAL modifiers are a PostgreSQL extension, as is the RESET syntax.

The privileges necessary to execute this command are left implementation-defined by the standard.

See Also

SET ROLE
SET TRANSACTION

Name

SET TRANSACTION — set the characteristics of the current transaction

Synopsis

SET TRANSACTION transaction_mode [, ...]
SET TRANSACTION SNAPSHOT snapshot_id
SET SESSION CHARACTERISTICS AS TRANSACTION transaction_mode [, ...]

where transaction_mode is one of:

   ISOLATION LEVEL | SERIALIZABLE | REPEATABLE READ | READ COMMITTED | READ UNCOMMITTED
   READ WRITE | READ ONLY
   [ NOT ] DEFERRABLE

Description

The SET TRANSACTION command sets the characteristics of the current transaction. It has no effect on any subsequent transactions. SET SESSION CHARACTERISTICS sets the default transaction characteristics for subsequent transactions of a session. These defaults can be overridden by SET TRANSACTION for an individual transaction.

The available transaction characteristics are the transaction isolation level, the transaction access mode (read/write or read-only), and the deferrable mode. In addition, a snapshot can be selected, though only for the current transaction, not as a session default.

The isolation level of a transaction determines what data the transaction can see when other transactions are running concurrently:

READ COMMITTED
   A statement can only see rows committed before it began. This is the default.

REPEATABLE READ
   All statements of the current transaction can only see rows committed before the first query or data-modification statement was executed in this transaction.

SERIALIZABLE
   All statements of the current transaction can only see rows committed before the first query or data-modification statement was executed in this transaction. If a pattern of reads and writes among concurrent serializable transactions would create a situation which could not have occurred for any serial (one-at-a-time) execution of those transactions, one of them will be rolled back with a serialization_failure error.

The SQL standard defines one additional level, READ UNCOMMITTED. In PostgreSQL READ UNCOMMITTED is treated as READ COMMITTED.

The transaction isolation level cannot be changed after the first query or data-modification statement (SELECT, INSERT, DELETE, UPDATE, FETCH, or COPY) of a transaction has been executed. See Chapter 13 for more information about transaction isolation and concurrency control.
The transaction access mode determines whether the transaction is read/write or read-only. Read/write is the default. When a transaction is read-only, the following SQL commands are disallowed: `INSERT`, `UPDATE`, `DELETE`, and `COPY FROM` if the table they would write to is not a temporary table; all `CREATE`, `ALTER`, and `DROP` commands; `COMMENT`, `GRANT`, `REVOKE`, `TRUNCATE`; and `EXPLAIN ANALYZE` and `EXECUTE` if the command they would execute is among those listed. This is a high-level notion of read-only that does not prevent all writes to disk.

The `DEFERRABLE` transaction property has no effect unless the transaction is also `SERIALIZABLE` and `READ ONLY`. When all three of these properties are selected for a transaction, the transaction may block when first acquiring its snapshot, after which it is able to run without the normal overhead of a `SERIALIZABLE` transaction and without any risk of contributing to or being canceled by a serialization failure. This mode is well suited for long-running reports or backups.

The `SET TRANSACTION SNAPSHOT` command allows a new transaction to run with the same snapshot as an existing transaction. The pre-existing transaction must have exported its snapshot with the `pg_export_snapshot` function (see Section 9.26.5). That function returns a snapshot identifier, which must be given to `SET TRANSACTION SNAPSHOT` to specify which snapshot is to be imported. The identifier must be written as a string literal in this command, for example `’000003A1-1’`. `SET TRANSACTION SNAPSHOT` can only be executed at the start of a transaction, before the first query or data-modification statement (`SELECT`, `INSERT`, `DELETE`, `UPDATE`, `FETCH`, or `COPY`) of the transaction. Furthermore, the transaction must already be set to `SERIALIZABLE` or `REPEATABLE READ` isolation level (otherwise, the snapshot would be discarded immediately, since `READ COMMITTED` mode takes a new snapshot for each command). If the importing transaction uses `SERIALIZABLE` isolation level, then the transaction that exported the snapshot must also use that isolation level. Also, a non-read-only serializable transaction cannot import a snapshot from a read-only transaction.

**Notes**

If `SET TRANSACTION` is executed without a prior `START TRANSACTION` or `BEGIN`, it emits a warning and otherwise has no effect.

It is possible to dispense with `SET TRANSACTION` by instead specifying the desired `transaction_modes` in `BEGIN` or `START TRANSACTION`. But that option is not available for `SET TRANSACTION SNAPSHOT`.

The session default transaction modes can also be set by setting the configuration parameters `default_transaction_isolation`, `default_transaction_read_only`, and `default_transaction_deferrable`. (In fact `SET SESSION CHARACTERISTICS` is just a verbose equivalent for setting these variables with `SET`.) This means the defaults can be set in the configuration file, via `ALTER DATABASE`, etc. Consult Chapter 19 for more information.

**Examples**

To begin a new transaction with the same snapshot as an already existing transaction, first export the snapshot from the existing transaction. That will return the snapshot identifier, for example:

```sql
BEGIN TRANSACTION ISOLATION LEVEL REPEATABLE READ;
SELECT pg_export_snapshot();
pg_export_snapshot
----------------------
000003A1-1
(1 row)
```
Then give the snapshot identifier in a SET TRANSACTION SNAPSHOT command at the beginning of the newly opened transaction:

```
BEGIN TRANSACTION ISOLATION LEVEL REPEATABLE READ;
SET TRANSACTION SNAPSHOT '000003A1-1';
```

**Compatibility**

These commands are defined in the SQL standard, except for the DEFERRABLE transaction mode and the SET TRANSACTION SNAPSHOT form, which are PostgreSQL extensions.

`SERIALIZABLE` is the default transaction isolation level in the standard. In PostgreSQL the default is ordinarily `READ COMMITTED`, but you can change it as mentioned above.

In the SQL standard, there is one other transaction characteristic that can be set with these commands: the size of the diagnostics area. This concept is specific to embedded SQL, and therefore is not implemented in the PostgreSQL server.

The SQL standard requires commas between successive `transaction_modes`, but for historical reasons PostgreSQL allows the commas to be omitted.
**SHOW**

**Name**

SHOW — show the value of a run-time parameter

**Synopsis**

SHOW name  
SHOW ALL

**Description**

SHOW will display the current setting of run-time parameters. These variables can be set using the SET statement, by editing the *postgresql.conf* configuration file, through the `PGOPTIONS` environmental variable (when using libpq or a libpq-based application), or through command-line flags when starting the *postgres* server. See Chapter 19 for details.

**Parameters**

*name*

The name of a run-time parameter. Available parameters are documented in Chapter 19 and on the SET reference page. In addition, there are a few parameters that can be shown but not set:

- **SERVER_VERSION**
  Shows the server’s version number.

- **SERVER_ENCODING**
  Shows the server-side character set encoding. At present, this parameter can be shown but not set, because the encoding is determined at database creation time.

- **LC_COLLATE**
  Shows the database’s locale setting for collation (text ordering). At present, this parameter can be shown but not set, because the setting is determined at database creation time.

- **LC_CTYPE**
  Shows the database’s locale setting for character classification. At present, this parameter can be shown but not set, because the setting is determined at database creation time.

- **IS_SUPERUSER**
  True if the current role has superuser privileges.

- **ALL**
  Show the values of all configuration parameters, with descriptions.
Notes

The function `current_setting` produces equivalent output; see Section 9.26. Also, the `pg_settings` system view produces the same information.

Examples

Show the current setting of the parameter `DateStyle`:

```
SHOW DateStyle;
DateStyle
-----------
ISO, MDY
```

Show the current setting of the parameter `geqo`:

```
SHOW geqo;
geqo
------
on
```

Show all settings:

```
SHOW ALL;

<table>
<thead>
<tr>
<th>name</th>
<th>setting</th>
<th>description</th>
</tr>
</thead>
<tbody>
<tr>
<td>allow_system_table_mods</td>
<td>off</td>
<td>Allows modifications of the structure of ...</td>
</tr>
<tr>
<td></td>
<td></td>
<td>.</td>
</tr>
<tr>
<td></td>
<td></td>
<td>.</td>
</tr>
<tr>
<td>xmloption</td>
<td>content</td>
<td>Sets whether XML data in implicit parsing ...</td>
</tr>
<tr>
<td>zero_damaged_pages</td>
<td>off</td>
<td>Continues processing past damaged page headers.</td>
</tr>
</tbody>
</table>
```

Compatibility

The `SHOW` command is a PostgreSQL extension.

See Also

`SET`, `RESET`
START TRANSACTION

Name
START TRANSACTION — start a transaction block

Synopsis
START TRANSACTION [ transaction_mode [, ...] ]

where transaction_mode is one of:

- ISOLATION LEVEL | SERIALIZABLE | REPEATABLE READ | READ COMMITTED | READ UNCOMMITTED
- READ WRITE | READ ONLY
- [ NOT ] DEFERRABLE

Description
This command begins a new transaction block. If the isolation level, read/write mode, or deferrable mode is specified, the new transaction has those characteristics, as if SET TRANSACTION was executed. This is the same as the BEGIN command.

Parameters
Refer to SET TRANSACTION for information on the meaning of the parameters to this statement.

Compatibility
In the standard, it is not necessary to issue START TRANSACTION to start a transaction block: any SQL command implicitly begins a block. PostgreSQL’s behavior can be seen as implicitly issuing a COMMIT after each command that does not follow START TRANSACTION (or BEGIN), and it is therefore often called “autocommit”. Other relational database systems might offer an autocommit feature as a convenience.

The DEFERRABLE transaction_mode is a PostgreSQL language extension.

The SQL standard requires commas between successive transaction_modes, but for historical reasons PostgreSQL allows the commas to be omitted.

See also the compatibility section of SET TRANSACTION.

See Also
BEGIN, COMMIT, ROLLBACK, SAVEPOINT, SET TRANSACTION
TRUNCATE

Name

TRUNCATE — empty a table or set of tables

Synopsis

TRUNCATE [ TABLE ] [ ONLY ] name [ * ] [, ... ]

[ RESTART IDENTITY | CONTINUE IDENTITY ] [ CASCADE | RESTRICT ]

Description

TRUNCATE quickly removes all rows from a set of tables. It has the same effect as an unqualified DELETE on each table, but since it does not actually scan the tables it is faster. Furthermore, it reclaims disk space immediately, rather than requiring a subsequent VACUUM operation. This is most useful on large tables.

Parameters

name

The name (optionally schema-qualified) of a table to truncate. If ONLY is specified before the table name, only that table is truncated. If ONLY is not specified, the table and all its descendant tables (if any) are truncated. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.

RESTART IDENTITY

Automatically restart sequences owned by columns of the truncated table(s).

CONTINUE IDENTITY

Do not change the values of sequences. This is the default.

CASCADE

Automatically truncate all tables that have foreign-key references to any of the named tables, or to any tables added to the group due to CASCADE.

RESTRICT

Refuse to truncate if any of the tables have foreign-key references from tables that are not listed in the command. This is the default.

Notes

You must have the TRUNCATE privilege on a table to truncate it.

TRUNCATE acquires an ACCESS EXCLUSIVE lock on each table it operates on, which blocks all other concurrent operations on the table. When RESTART IDENTITY is specified, any sequences that are to be restarted are likewise locked exclusively. If concurrent access to a table is required, then the DELETE command should be used instead.
TRUNCATE cannot be used on a table that has foreign-key references from other tables, unless all such tables are also truncated in the same command. Checking validity in such cases would require table scans, and the whole point is not to do one. The CASCADE option can be used to automatically include all dependent tables — but be very careful when using this option, or else you might lose data you did not intend to!

TRUNCATE will not fire any ON DELETE triggers that might exist for the tables. But it will fire ON TRUNCATE triggers. If ON TRUNCATE triggers are defined for any of the tables, then all BEFORE TRUNCATE triggers are fired before any truncation happens, and all AFTER TRUNCATE triggers are fired after the last truncation is performed and any sequences are reset. The triggers will fire in the order that the tables are to be processed (first those listed in the command, and then any that were added due to cascading).

TRUNCATE is not MVCC-safe. After truncation, the table will appear empty to concurrent transactions, if they are using a snapshot taken before the truncation occurred. See Section 13.5 for more details. TRUNCATE is transaction-safe with respect to the data in the tables: the truncation will be safely rolled back if the surrounding transaction does not commit.

When RESTART IDENTITY is specified, the implied ALTER SEQUENCE RESTART operations are also done transactionally; that is, they will be rolled back if the surrounding transaction does not commit. This is unlike the normal behavior of ALTER SEQUENCE RESTART. Be aware that if any additional sequence operations are done on the restarted sequences before the transaction rolls back, the effects of these operations on the sequences will be rolled back, but not their effects on currval(); that is, after the transaction currval() will continue to reflect the last sequence value obtained inside the failed transaction, even though the sequence itself may no longer be consistent with that. This is similar to the usual behavior of currval() after a failed transaction.

TRUNCATE is not currently supported for foreign tables. This implies that if a specified table has any descendant tables that are foreign, the command will fail.

Examples

Truncate the tables bigtable and fattable:

TRUNCATE bigtable, fattable;

The same, and also reset any associated sequence generators:

TRUNCATE bigtable, fattable RESTART IDENTITY;

Truncate the table othertable, and cascade to any tables that reference othertable via foreign-key constraints:

TRUNCATE othertable CASCADE;

Compatibility

The SQL:2008 standard includes a TRUNCATE command with the syntax TRUNCATE TABLE tablename. The clauses CONTINUE IDENTITY/RESTART IDENTITY also appear in that standard,
but have slightly different though related meanings. Some of the concurrency behavior of this command is left implementation-defined by the standard, so the above notes should be considered and compared with other implementations if necessary.
**UNLISTEN**

**Name**
UNLISTEN — stop listening for a notification

**Synopsis**
UNLISTEN { channel | * }

**Description**
UNLISTEN is used to remove an existing registration for NOTIFY events. UNLISTEN cancels any existing registration of the current PostgreSQL session as a listener on the notification channel named `channel`. The special wildcard `*` cancels all listener registrations for the current session.

NOTIFY contains a more extensive discussion of the use of LISTEN and NOTIFY.

**Parameters**
- **channel**
  - Name of a notification channel (any identifier).
- *****
  - All current listen registrations for this session are cleared.

**Notes**
You can unlisten something you were not listening for; no warning or error will appear.

At the end of each session, UNLISTEN * is automatically executed.

A transaction that has executed UNLISTEN cannot be prepared for two-phase commit.

**Examples**
To make a registration:

LISTEN virtual;
NOTIFY virtual;
Asynchronous notification "virtual" received from server process with PID 8448.

Once UNLISTEN has been executed, further NOTIFY messages will be ignored:

UNLISTEN virtual;
NOTIFY virtual;
-- no NOTIFY event is received
Compatibility

There is no UNLISTEN command in the SQL standard.

See Also

LISTEN, NOTIFY
UPDATE

Name

UPDATE — update rows of a table

Synopsis

[ WITH [ RECURSIVE ] with_query [, ... ] ]
UPDATE [ ONLY ] table_name [ * ] [ [ AS ] alias ]
SET { column_name = { expression | DEFAULT } |
( column_name [, ...] ) = ( { expression | DEFAULT } [, ...] ) |
( column_name [, ...] ) = ( sub-SELECT )
} [, ...]
[ FROM from_list ]
[ WHERE condition | WHERE CURRENT OF cursor_name ]
[ RETURNING * | output_expression [ [ AS ] output_name ] [, ...] ]

Description

UPDATE changes the values of the specified columns in all rows that satisfy the condition. Only the columns to be modified need be mentioned in the SET clause; columns not explicitly modified retain their previous values.

There are two ways to modify a table using information contained in other tables in the database: using sub-selects, or specifying additional tables in the FROM clause. Which technique is more appropriate depends on the specific circumstances.

The optional RETURNING clause causes UPDATE to compute and return value(s) based on each row actually updated. Any expression using the table’s columns, and/or columns of other tables mentioned in FROM, can be computed. The new (post-update) values of the table’s columns are used. The syntax of the RETURNING list is identical to that of the output list of SELECT.

You must have the UPDATE privilege on the table, or at least on the column(s) that are listed to be updated. You must also have the SELECT privilege on any column whose values are read in the expressions or condition.

Parameters

with_query

The WITH clause allows you to specify one or more subqueries that can be referenced by name in the UPDATE query. See Section 7.8 and SELECT for details.

table_name

The name (optionally schema-qualified) of the table to update. If ONLY is specified before the table name, matching rows are updated in the named table only. If ONLY is not specified, matching rows are also updated in any tables inheriting from the named table. Optionally, * can be specified after the table name to explicitly indicate that descendant tables are included.
**UPDATE**

**alias**

A substitute name for the target table. When an alias is provided, it completely hides the actual name of the table. For example, given `UPDATE foo AS f`, the remainder of the `UPDATE` statement must refer to this table as `f` not `foo`.

**column_name**

The name of a column in the table named by `table_name`. The column name can be qualified with a subfield name or array subscript, if needed. Do not include the table’s name in the specification of a target column — for example, `UPDATE table_name SET table_name.col = 1` is invalid.

**expression**

An expression to assign to the column. The expression can use the old values of this and other columns in the table.

**DEFAULT**

Set the column to its default value (which will be NULL if no specific default expression has been assigned to it).

**sub-SELECT**

A `SELECT` sub-query that produces as many output columns as are listed in the parenthesized column list preceding it. The sub-query must yield no more than one row when executed. If it yields one row, its column values are assigned to the target columns; if it yields no rows, NULL values are assigned to the target columns. The sub-query can refer to old values of the current row of the table being updated.

**from_list**

A list of table expressions, allowing columns from other tables to appear in the `WHERE` condition and the update expressions. This is similar to the list of tables that can be specified in the `FROM Clause` of a `SELECT` statement. Note that the target table must not appear in the `from_list`, unless you intend a self-join (in which case it must appear with an alias in the `from_list`).

**condition**

An expression that returns a value of type `boolean`. Only rows for which this expression returns `true` will be updated.

**cursor_name**

The name of the cursor to use in a `WHERE CURRENT OF` condition. The row to be updated is the one most recently fetched from this cursor. The cursor must be a non-grouping query on the `UPDATE`’s target table. Note that `WHERE CURRENT OF` cannot be specified together with a Boolean condition. See DECLARE for more information about using cursors with `WHERE CURRENT OF`.

**output_expression**

An expression to be computed and returned by the `UPDATE` command after each row is updated. The expression can use any column names of the table named by `table_name` or table(s) listed in `FROM`. Write `-` to return all columns.

**output_name**

A name to use for a returned column.
**Outputs**

On successful completion, an `UPDATE` command returns a command tag of the form

```
UPDATE count
```

The `count` is the number of rows updated, including matched rows whose values did not change. Note that the number may be less than the number of rows that matched the `condition` when updates were suppressed by a `BEFORE UPDATE` trigger. If `count` is 0, no rows were updated by the query (this is not considered an error).

If the `UPDATE` command contains a `RETURNING` clause, the result will be similar to that of a `SELECT` statement containing the columns and values defined in the `RETURNING` list, computed over the row(s) updated by the command.

**Notes**

When a `FROM` clause is present, what essentially happens is that the target table is joined to the tables mentioned in the `from_list`, and each output row of the join represents an update operation for the target table. When using `FROM` you should ensure that the join produces at most one output row for each row to be modified. In other words, a target row shouldn’t join to more than one row from the other table(s). If it does, then only one of the join rows will be used to update the target row, but which one will be used is not readily predictable.

Because of this indeterminacy, referencing other tables only within sub-selects is safer, though often harder to read and slower than using a join.

**Examples**

Change the word `Drama` to `Dramatic` in the column `kind` of the table `films`:

```
UPDATE films SET kind = 'Dramatic' WHERE kind = 'Drama';
```

Adjust temperature entries and reset precipitation to its default value in one row of the table `weather`:

```
UPDATE weather SET temp_lo = temp_lo+1, temp_hi = temp_lo+15, prcp = DEFAULT
WHERE city = 'San Francisco' AND date = '2003-07-03';
```

Perform the same operation and return the updated entries:

```
UPDATE weather SET temp_lo = temp_lo+1, temp_hi = temp_lo+15, prcp = DEFAULT
WHERE city = 'San Francisco' AND date = '2003-07-03'
RETURNING temp_lo, temp_hi, prcp;
```

Use the alternative column-list syntax to do the same update:

```
UPDATE weather SET (temp_lo, temp_hi, prcp) = (temp_lo+1, temp_lo+15, DEFAULT)
WHERE city = 'San Francisco' AND date = '2003-07-03';
```
Increment the sales count of the salesperson who manages the account for Acme Corporation, using the **FROM** clause syntax:

```
UPDATE employees SET sales_count = sales_count + 1 FROM accounts
    WHERE accounts.name = 'Acme Corporation'
    AND employees.id = accounts.sales_person;
```

Perform the same operation, using a sub-select in the **WHERE** clause:

```
UPDATE employees SET sales_count = sales_count + 1 WHERE id =
    (SELECT sales_person FROM accounts WHERE name = 'Acme Corporation');
```

Update contact names in an accounts table to match the currently assigned salesmen:

```
UPDATE accounts SET (contact_first_name, contact_last_name) =
    (SELECT first_name, last_name FROM salesmen
        WHERE salesmen.id = accounts.sales_id);
```

A similar result could be accomplished with a join:

```
UPDATE accounts SET contact_first_name = first_name,
    contact_last_name = last_name
    FROM salesmen WHERE salesmen.id = accounts.sales_id;
```

However, the second query may give unexpected results if `salesmen.id` is not a unique key, whereas the first query is guaranteed to raise an error if there are multiple `id` matches. Also, if there is no match for a particular `accounts.sales_id` entry, the first query will set the corresponding name fields to NULL, whereas the second query will not update that row at all.

Update statistics in a summary table to match the current data:

```
UPDATE summary s SET (sum_x, sum_y, avg_x, avg_y) =
    (SELECT sum(x), sum(y), avg(x), avg(y) FROM data d
        WHERE d.group_id = s.group_id);
```

Attempt to insert a new stock item along with the quantity of stock. If the item already exists, instead update the stock count of the existing item. To do this without failing the entire transaction, use savepoints:

```
BEGIN;
    -- other operations
    SAVEPOINT sp1;
    INSERT INTO wines VALUES('Chateau Lafite 2003', '24');
    -- Assume the above fails because of a unique key violation,
    -- so now we issue these commands:
    ROLLBACK TO sp1;
    UPDATE wines SET stock = stock + 24 WHERE winename = 'Chateau Lafite 2003';
    -- continue with other operations, and eventually
    COMMIT;
```

Change the `kind` column of the table `films` in the row on which the cursor `c_films` is currently positioned:
UPDATE films SET kind = 'Dramatic' WHERE CURRENT OF c_films;

Compatibility

This command conforms to the SQL standard, except that the FROM and RETURNING clauses are PostgreSQL extensions, as is the ability to use WITH with UPDATE.

Some other database systems offer a FROM option in which the target table is supposed to be listed again within FROM. That is not how PostgreSQL interprets FROM. Be careful when porting applications that use this extension.

According to the standard, the source value for a parenthesized sub-list of column names can be any row-valued expression yielding the correct number of columns. PostgreSQL only allows the source value to be a parenthesized list of expressions or a sub-SELECT. An individual column’s updated value can be specified as DEFAULT in the list-of-expressions case, but not inside a sub-SELECT.
VACUUM

Name

VACUUM — garbage-collect and optionally analyze a database

Synopsis

VACUUM [ ( { FULL | FREEZE | VERBOSE | ANALYZE | DISABLE_PAGE_SKIPPING } [, ...] ) ] [ table_name ]
VACUUM [ FULL ] [ FREEZE ] [ VERBOSE ] [ table_name ]
VACUUM [ FULL ] [ FREEZE ] [ VERBOSE ] ANALYZE [ table_name [ (column_name [, ...] ) ] ]

Description

VACUUM reclaims storage occupied by dead tuples. In normal PostgreSQL operation, tuples that are deleted or obsoleted by an update are not physically removed from their table; they remain present until a VACUUM is done. Therefore it’s necessary to do VACUUM periodically, especially on frequently-updated tables.

With no parameter, VACUUM processes every table in the current database that the current user has permission to vacuum. With a parameter, VACUUM processes only that table.

VACUUM ANALYZE performs a VACUUM and then an ANALYZE for each selected table. This is a handy combination form for routine maintenance scripts. See ANALYZE for more details about its processing.

Plain VACUUM (without FULL) simply reclaims space and makes it available for re-use. This form of the command can operate in parallel with normal reading and writing of the table, as an exclusive lock is not obtained. However, extra space is not returned to the operating system (in most cases); it’s just kept available for re-use within the same table. VACUUM FULL rewrites the entire contents of the table into a new disk file with no extra space, allowing unused space to be returned to the operating system. This form is much slower and requires an exclusive lock on each table while it is being processed.

When the option list is surrounded by parentheses, the options can be written in any order. Without parentheses, options must be specified in exactly the order shown above. The parenthesized syntax was added in PostgreSQL 9.0; the unparenthesized syntax is deprecated.

Parameters

FULL

Selects “full” vacuum, which can reclaim more space, but takes much longer and exclusively locks the table. This method also requires extra disk space, since it writes a new copy of the table and doesn’t release the old copy until the operation is complete. Usually this should only be used when a significant amount of space needs to be reclaimed from within the table.

FREEZE

Selects aggressive “freezing” of tuples. Specifying FREEZE is equivalent to performing VACUUM with the vacuum_freeze_min_age and vacuum_freeze_table_age parameters set to zero. Aggressive freezing is always performed when the table is rewritten, so this option is redundant when FULL is specified.
VACUUM

VERBOSE

Prints a detailed vacuum activity report for each table.

ANALYZE

Updates statistics used by the planner to determine the most efficient way to execute a query.

DISABLE_PAGE_SKIPPING

Normally, VACUUM will skip pages based on the visibility map. Pages where all tuples are known to be frozen can always be skipped, and those where all tuples are known to be visible to all transactions may be skipped except when performing an aggressive vacuum. Furthermore, except when performing an aggressive vacuum, some pages may be skipped in order to avoid waiting for other sessions to finish using them. This option disables all page-skipping behavior, and is intended to be used only the contents of the visibility map are thought to be suspect, which should happen only if there is a hardware or software issue causing database corruption.

table_name

The name (optionally schema-qualified) of a specific table to vacuum. Defaults to all tables in the current database.

column_name

The name of a specific column to analyze. Defaults to all columns. If a column list is specified, ANALYZE is implied.

Outputs

When VERBOSE is specified, VACUUM emits progress messages to indicate which table is currently being processed. Various statistics about the tables are printed as well.

Notes

To vacuum a table, one must ordinarily be the table’s owner or a superuser. However, database owners are allowed to vacuum all tables in their databases, except shared catalogs. (The restriction for shared catalogs means that a true database-wide VACUUM can only be performed by a superuser.) VACUUM will skip over any tables that the calling user does not have permission to vacuum.

VACUUM cannot be executed inside a transaction block.

For tables with GIN indexes, VACUUM (in any form) also completes any pending index insertions, by moving pending index entries to the appropriate places in the main GIN index structure. See Section 63.4.1 for details.

We recommend that active production databases be vacuumed frequently (at least nightly), in order to remove dead rows. After adding or deleting a large number of rows, it might be a good idea to issue a VACUUM ANALYZE command for the affected table. This will update the system catalogs with the results of all recent changes, and allow the PostgreSQL query planner to make better choices in planning queries.

The FULL option is not recommended for routine use, but might be useful in special cases. An example is when you have deleted or updated most of the rows in a table and would like the table to physically shrink to occupy less disk space and allow faster table scans. VACUUM FULL will usually shrink the table more than a plain VACUUM would.
VACUUM causes a substantial increase in I/O traffic, which might cause poor performance for other active sessions. Therefore, it is sometimes advisable to use the cost-based vacuum delay feature. See Section 19.4.4 for details.

PostgreSQL includes an “autovacuum” facility which can automate routine vacuum maintenance. For more information about automatic and manual vacuuming, see Section 24.1.

Examples
To clean a single table onek, analyze it for the optimizer and print a detailed vacuum activity report:

```
VACUUM (VERBOSE, ANALYZE) onek;
```

Compatibility
There is no VACUUM statement in the SQL standard.

See Also
vacuumdb, Section 19.4.4, Section 24.1.6
VALUES

Name
VALUES — compute a set of rows

Synopsis
VALUES ( expression [, ...] ) [, ...]
[ ORDER BY sort_expression [ ASC | DESC | USING operator ] [, ... ] ]
[ LIMIT { count | ALL } ]
[ OFFSET start [ ROW | ROWS ] ]
[ FETCH { FIRST | NEXT } [ count ] { ROW | ROWS } ONLY ]

Description
VALUES computes a row value or set of row values specified by value expressions. It is most com-
monly used to generate a “constant table” within a larger command, but it can be used on its own.

When more than one row is specified, all the rows must have the same number of elements. The data
types of the resulting table’s columns are determined by combining the explicit or inferred types of
the expressions appearing in that column, using the same rules as for UNION (see Section 10.5).

Within larger commands, VALUES is syntactically allowed anywhere that SELECT is. Because it is
treated like a SELECT by the grammar, it is possible to use the ORDER BY, LIMIT (or equivalently
FETCH FIRST), and OFFSET clauses with a VALUES command.

Parameters

expression
A constant or expression to compute and insert at the indicated place in the resulting table (set of
rows). In a VALUES list appearing at the top level of an INSERT, an expression can be replaced
by DEFAULT to indicate that the destination column’s default value should be inserted. DEFAULT
cannot be used when VALUES appears in other contexts.

sort_expression
An expression or integer constant indicating how to sort the result rows. This expression can refer
to the columns of the VALUES result as column1, column2, etc. For more details see ORDER
BY Clause.

operator
A sorting operator. For details see ORDER BY Clause.

count
The maximum number of rows to return. For details see LIMIT Clause.

start
The number of rows to skip before starting to return rows. For details see LIMIT Clause.
Notes

VALUES lists with very large numbers of rows should be avoided, as you might encounter out-of-memory failures or poor performance. VALUES appearing within INSERT is a special case (because the desired column types are known from the INSERT's target table, and need not be inferred by scanning the VALUES list), so it can handle larger lists than are practical in other contexts.

Examples

A bare VALUES command:

VALUES (1, 'one'), (2, 'two'), (3, 'three');

This will return a table of two columns and three rows. It’s effectively equivalent to:

SELECT 1 AS column1, 'one' AS column2
UNION ALL
SELECT 2, 'two'
UNION ALL
SELECT 3, 'three';

More usually, VALUES is used within a larger SQL command. The most common use is in INSERT:

INSERT INTO films (code, title, did, date_prod, kind)
VALUES ('T_601', 'Yojimbo', 106, '1961-06-16', 'Drama');

In the context of INSERT, entries of a VALUES list can be DEFAULT to indicate that the column default should be used here instead of specifying a value:

INSERT INTO films VALUES
    ('UA502', 'Bananas', 105, DEFAULT, 'Comedy', '82 minutes'),
    ('T_601', 'Yojimbo', 106, DEFAULT, 'Drama', DEFAULT);

VALUES can also be used where a sub-SELECT might be written, for example in a FROM clause:

SELECT f.*
FROM films f, (VALUES('MGM', 'Horror'), ('UA', 'Sci-Fi')) AS t (studio, kind)
WHERE f.studio = t.studio AND f.kind = t.kind;

UPDATE employees SET salary = salary * v.increase
FROM (VALUES(1, 200000, 1.2), (2, 400000, 1.4)) AS v (depno, target, increase)
WHERE employees.depno = v.depno AND employees.sales >= v.target;

Note that an AS clause is required when VALUES is used in a FROM clause, just as is true for SELECT. It is not required that the AS clause specify names for all the columns, but it’s good practice to do so. (The default column names for VALUES are column1, column2, etc in PostgreSQL, but these names might be different in other database systems.)

When VALUES is used in INSERT, the values are all automatically coerced to the data type of the corresponding destination column. When it’s used in other contexts, it might be necessary to specify the correct data type. If the entries are all quoted literal constants, coercing the first is sufficient to determine the assumed type for all:
VALUES

SELECT * FROM machines
WHERE ip_address IN (VALUES('192.168.0.1'::inet), ('192.168.0.10'), ('192.168.1.43'));

Tip: For simple IN tests, it's better to rely on the list-of-scalars form of IN than to write a VALUES query as shown above. The list of scalars method requires less writing and is often more efficient.

Compatibility

VALUES conforms to the SQL standard. LIMIT and OFFSET are PostgreSQL extensions; see also under SELECT.

See Also

INSERT, SELECT
II. PostgreSQL Client Applications

This part contains reference information for PostgreSQL client applications and utilities. Not all of these commands are of general utility; some might require special privileges. The common feature of these applications is that they can be run on any host, independent of where the database server resides.

When specified on the command line, user and database names have their case preserved — the presence of spaces or special characters might require quoting. Table names and other identifiers do not have their case preserved, except where documented, and might require quoting.
clusterdb

Name

clusterdb — cluster a PostgreSQL database

Synopsis

clusterdb [connection-option...] [--verbose|-v] [--table|-t table] ... [dbname]

clusterdb [connection-option...] [--verbose|-v] --all|--all

Description

clusterdb is a utility for reclustering tables in a PostgreSQL database. It finds tables that have previously been clustered, and clusters them again on the same index that was last used. Tables that have never been clustered are not affected.

clusterdb is a wrapper around the SQL command CLUSTER. There is no effective difference between clustering databases via this utility and via other methods for accessing the server.

Options

clusterdb accepts the following command-line arguments:

- **-a**
  --all
  Cluster all databases.

- **-d** dbname
  --dbname= dbname
  Specifies the name of the database to be clustered. If this is not specified and -a (or --all) is not used, the database name is read from the environment variable PGDATABASE. If that is not set, the user name specified for the connection is used.

- **-e**
  --echo
  Echo the commands that clusterdb generates and sends to the server.

- **-q**
  --quiet
  Do not display progress messages.
-t table
--table=table

Cluster table only. Multiple tables can be clustered by writing multiple -t switches.

-v
--verbose

Print detailed information during processing.

-V
--version

Print the clusterdb version and exit.

-?
--help

Show help about clusterdb command line arguments, and exit.

clusterdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force clusterdb to prompt for a password before connecting to a database.

This option is never essential, since clusterdb will automatically prompt for a password if the server demands password authentication. However, clusterdb will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

--maintenance-db=dbname

Specifies the name of the database to connect to discover what other databases should be clustered. If not specified, the postgres database will be used, and if that does not exist, template1 will be used.
Environment

PGDATABASE
PGHOST
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).

Diagnostics

In case of difficulty, see CLUSTER and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Examples

To cluster the database test:

$ clusterdb test

To cluster a single table foo in a database named xyzzy:

$ clusterdb --table foo xyzzy

See Also

CLUSTER
createdb

Name
createdb — create a new PostgreSQL database

Synopsis
createdb [connection-option...] [option...] [dbname [description]]

Description
createdb creates a new PostgreSQL database.

Normally, the database user who executes this command becomes the owner of the new database. However, a different owner can be specified via the -O option, if the executing user has appropriate privileges.

createdb is a wrapper around the SQL command CREATE DATABASE. There is no effective difference between creating databases via this utility and via other methods for accessing the server.

Options
createdb accepts the following command-line arguments:

dbname
  Specifies the name of the database to be created. The name must be unique among all PostgreSQL databases in this cluster. The default is to create a database with the same name as the current system user.

description
  Specifies a comment to be associated with the newly created database.

-D tablespace
--tablespace=tablespace
  Specifies the default tablespace for the database. (This name is processed as a double-quoted identifier.)

-e
--echo
  Echo the commands that createdb generates and sends to the server.

-E encoding
--encoding=encoding
  Specifies the character encoding scheme to be used in this database. The character sets supported by the PostgreSQL server are described in Section 23.3.1.
-l locale
--locale=locale

Specifies the locale to be used in this database. This is equivalent to specifying both
--lc-collate and --lc-ctype.

--lc-collate=locale

Specifies the LC_COLLATE setting to be used in this database.

--lc-ctype=locale

Specifies the LC_CTYPE setting to be used in this database.

-O owner
--owner=owner

Specifies the database user who will own the new database. (This name is processed as a double-quoted identifier.)

-T template
--template=template

Specifies the template database from which to build this database. (This name is processed as a double-quoted identifier.)

-V
--version

Print the createdb version and exit.

-?
--help

Show help about createdb command line arguments, and exit.

The options -D, -l, -E, -O, and -T correspond to options of the underlying SQL command CREATE DATABASE; see there for more information about them.

createdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or the local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.
-W
--password

Force createdb to prompt for a password before connecting to a database.
This option is never essential, since createdb will automatically prompt for a password if the
server demands password authentication. However, createdb will waste a connection attempt
finding out that the server wants a password. In some cases it is worth typing -W to avoid the
extra connection attempt.

--maintenance-db=dbname

Specifies the name of the database to connect to when creating the new database. If not specified,
the postgres database will be used; if that does not exist (or if it is the name of the new database
being created), template1 will be used.

Environment

PGDATABASE

If set, the name of the database to create, unless overridden on the command line.

PGHOST
PGPORT
PGUSER

Default connection parameters. PGUSER also determines the name of the database to create, if it
is not specified on the command line or by PGDATABASE.

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by
libpq (see Section 32.14).

Diagnostics

In case of difficulty, see CREATE DATABASE and psql for discussions of potential problems and
error messages. The database server must be running at the targeted host. Also, any default connection
settings and environment variables used by the libpq front-end library will apply.

Examples

To create the database demo using the default database server:

$ createdb demo

To create the database demo using the server on host eden, port 5000, using the LATIN1 encoding
scheme with a look at the underlying command:

$ createdb -p 5000 -h eden -E LATIN1 -e demo
CREATE DATABASE demo ENCODING 'LATIN1';
See Also

dropdb, CREATE DATABASE
createlang

Name
createlang — install a PostgreSQL procedural language

Synopsis
createlang [connection-option...] langname [dbname]

createlang [connection-option...] --list -l [dbname]

Description
createlang is a utility for adding a procedural language to a PostgreSQL database.
createlang is just a wrapper around the CREATE EXTENSION SQL command.

Caution
createlang is deprecated and may be removed in a future PostgreSQL release.
Direct use of the CREATE EXTENSION command is recommended instead.

Options
createlang accepts the following command-line arguments:

langname
Specifies the name of the procedural language to be installed. (This name is lower-cased.)

[-d] dbname
[--dbname=] dbname
Specifies the database to which the language should be added. The default is to use the database with the same name as the current system user.

-e
--echo
Display SQL commands as they are executed.

-l
--list
Show a list of already installed languages in the target database.

-V
--version
Print the createlang version and exit.
createlang

-?
--help

Show help about createlang command line arguments, and exit.

createlang also accepts the following command-line arguments for connection parameters:

-\ h  \ host
   --host=host

   Specifies the host name of the machine on which the server is running. If the value begins with a
   slash, it is used as the directory for the Unix domain socket.

-\ p  \ port
   --port=port

   Specifies the TCP port or local Unix domain socket file extension on which the server is listening
   for connections.

-\ U  \ username
   --username=username

   User name to connect as.

-\ w
   --no-password

   Never issue a password prompt. If the server requires password authentication and a password
   is not available by other means such as a .pgpass file, the connection attempt will fail. This
   option can be useful in batch jobs and scripts where no user is present to enter a password.

-\ W
   --password

   Force createlang to prompt for a password before connecting to a database.

   This option is never essential, since createlang will automatically prompt for a password if the
   server demands password authentication. However, createlang will waste a connection attempt
   finding out that the server wants a password. In some cases it is worth typing -h to avoid the
   extra connection attempt.

Environment

PGDATABASE
PGHOST
PGPORT
PGUSER

   Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by
libpq (see Section 32.14).
Diagnostics

Most error messages are self-explanatory. If not, run createlang with the \-echo option and see the respective SQL command for details. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Notes

Use droplang to remove a language.

Examples

To install the language pltcl into the database template1:

$ createlang pltcl template1

Note that installing the language into template1 will cause it to be automatically installed into subsequently-created databases as well.

See Also

droplang, CREATE EXTENSION, CREATE LANGUAGE
createuser

Name
createuser — define a new PostgreSQL user account

Synopsis
createuser [connection-option...][option...][username]

Description
createuser creates a new PostgreSQL user (or more precisely, a role). Only superusers and users with CREATEROLE privilege can create new users, so createuser must be invoked by someone who can connect as a superuser or a user with CREATEROLE privilege.

If you wish to create a new superuser, you must connect as a superuser, not merely with CREATEROLE privilege. Being a superuser implies the ability to bypass all access permission checks within the database, so superuserdom should not be granted lightly.

createuser is a wrapper around the SQL command CREATE ROLE. There is no effective difference between creating users via this utility and via other methods for accessing the server.

Options
createuser accepts the following command-line arguments:

username
  Specifies the name of the PostgreSQL user to be created. This name must be different from all existing roles in this PostgreSQL installation.

-c number
--connection-limit=number
  Set a maximum number of connections for the new user. The default is to set no limit.

-d
--createdb
  The new user will be allowed to create databases.

-D
--no-createdb
  The new user will not be allowed to create databases. This is the default.

-e
--echo
  Echo the commands that createuser generates and sends to the server.
createuser

-Encrypted

Encrypts the user’s password stored in the database. If not specified, the default password behavior is used.

-g role

Indicates role to which this role will be added immediately as a new member. Multiple roles to which this role will be added as a member can be specified by writing multiple -g switches.

-i

--inherit

The new role will automatically inherit privileges of roles it is a member of. This is the default.

-I

--no-inherit

The new role will not automatically inherit privileges of roles it is a member of.

--interactive

Prompt for the user name if none is specified on the command line, and also prompt for whichever of the options -d/-D, -r/-R, -s/-S is not specified on the command line. (This was the default behavior up to PostgreSQL 9.1.)

-l

--login

The new user will be allowed to log in (that is, the user name can be used as the initial session user identifier). This is the default.

-L

--no-login

The new user will not be allowed to log in. (A role without login privilege is still useful as a means of managing database permissions.)

-N

--unencrypted

Does not encrypt the user’s password stored in the database. If not specified, the default password behavior is used.

-P

--pwprompt

If given, createuser will issue a prompt for the password of the new user. This is not necessary if you do not plan on using password authentication.

-r

--createrole

The new user will be allowed to create new roles (that is, this user will have CREATEROLE privilege).

-R

--no-createrole

The new user will not be allowed to create new roles. This is the default.
createuser

-s
--superuser
The new user will be a superuser.

-S
--no-superuser
The new user will not be a superuser. This is the default.

-V
--version
Print the createuser version and exit.

--replication
The new user will have the REPLICATION privilege, which is described more fully in the documentation for CREATE ROLE.

--no-replication
The new user will not have the REPLICATION privilege, which is described more fully in the documentation for CREATE ROLE.

-?
--help
Show help about createuser command line arguments, and exit.

createuser also accepts the following command-line arguments for connection parameters:

-h host
--host=host
Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

-p port
--port=port
Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

-U username
--username=username
User name to connect as (not the user name to create).

-w
--no-password
Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password
Force createuser to prompt for a password (for connecting to the server, not for the password of the new user).

This option is never essential, since createuser will automatically prompt for a password if the server demands password authentication. However, createuser will waste a connection attempt

1778
finding out that the server wants a password. In some cases it is worth typing `-w` to avoid the extra connection attempt.

Environment

**PGHOST**
**PGPORT**
**PGUSER**

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).

Diagnostics

In case of difficulty, see CREATE ROLE and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Examples

To create a user `joe` on the default database server:

```
$ createuser joe
```

To create a user `joe` on the default database server with prompting for some additional attributes:

```
$ createuser --interactive joe
Shall the new role be a superuser? (y/n) n
Shall the new role be allowed to create databases? (y/n) n
Shall the new role be allowed to create more new roles? (y/n) n
```

To create the same user `joe` using the server on host `eden`, port 5000, with attributes explicitly specified, taking a look at the underlying command:

```
$ createuser -h eden -p 5000 -S -D -R -e joe
CREATE ROLE joe NOSUPERUSER NOCREATEDB NOCREATEROLE INHERIT LOGIN;
```

To create the user `joe` as a superuser, and assign a password immediately:

```
$ createuser -P -s -e joe
Enter password for new role: xyzzy
Enter it again: xyzzy
CREATE ROLE joe PASSWORD 'md5b5f5ba1a423792b526f799ae4eb3d59e' SUPERUSER CREATEDBCREATEROLE INHERIT LOGIN;
```
In the above example, the new password isn’t actually echoed when typed, but we show what was typed for clarity. As you see, the password is encrypted before it is sent to the client. If the option --unencrypted is used, the password will appear in the echoed command (and possibly also in the server log and elsewhere), so you don’t want to use --e in that case, if anyone else can see your screen.

**See Also**

dropuser, CREATE ROLE
**dropdb**

**Name**

*dropdb* — remove a PostgreSQL database

**Synopsis**

```
dropdb [connection-option...][option...] dbname
```

**Description**

*dropdb* destroys an existing PostgreSQL database. The user who executes this command must be a database superuser or the owner of the database.

*dropdb* is a wrapper around the SQL command DROP DATABASE. There is no effective difference between dropping databases via this utility and via other methods for accessing the server.

**Options**

dropdb accepts the following command-line arguments:

*dbname*

  Specifies the name of the database to be removed.

- **-e**  
  --echo

  Echo the commands that *dropdb* generates and sends to the server.

- **-i**  
  --interactive

  Issues a verification prompt before doing anything destructive.

- **-V**  
  --version

  Print the *dropdb* version and exit.

- **--if-exists**

  Do not throw an error if the database does not exist. A notice is issued in this case.

- **-?**  
  --help

  Show help about *dropdb* command line arguments, and exit.

*dropdb* also accepts the following command-line arguments for connection parameters:
-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a
slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening
for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password
is not available by other means such as a .pgpass file, the connection attempt will fail. This
option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force dropdb to prompt for a password before connecting to a database.

This option is never essential, since dropdb will automatically prompt for a password if the server
demands password authentication. However, dropdb will waste a connection attempt finding out
that the server wants a password. In some cases it is worth typing -W to avoid the extra connection
attempt.

--maintenance-db=dbname

Specifies the name of the database to connect to in order to drop the target database. If not
specified, the postgres database will be used; if that does not exist (or is the database being
dropped), template1 will be used.

Environment

PGHOST
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by
libpq (see Section 32.14).

Diagnostics

In case of difficulty, see DROP DATABASE and psql for discussions of potential problems and error
messages. The database server must be running at the targeted host. Also, any default connection
settings and environment variables used by the libpq front-end library will apply.
Examples

To destroy the database demo on the default database server:

$ dropdb demo

To destroy the database demo using the server on host eden, port 5000, with verification and a peek at the underlying command:

$ dropdb -p 5000 -h eden -i -e demo
Database "demo" will be permanently deleted.
Are you sure? (y/n) y
DROP DATABASE demo;

See Also

createdb, DROP DATABASE
droplang

Name
droplang — remove a PostgreSQL procedural language

Synopsis
droplang [connection-option...] langname [dbname]
droplang [connection-option...]|-l [dbname]

Description
droplang is a utility for removing an existing procedural language from a PostgreSQL database.
droplang is just a wrapper around the DROP EXTENSION SQL command.

Caution
droplang is deprecated and may be removed in a future PostgreSQL release.
Direct use of the DROP EXTENSION command is recommended instead.

Options
droplang accepts the following command line arguments:

langname
    Specifies the name of the procedural language to be removed. (This name is lower-cased.)

[-d] dbname
    Specifies from which database the language should be removed. The default is to use the database
    with the same name as the current system user.

-e
    --echo
    Display SQL commands as they are executed.

-1
    --list
    Show a list of already installed languages in the target database.

-V
    --version
    Print the droplang version and exit.
droplang

-?
--help

Show help about droplang command line arguments, and exit.

droplang also accepts the following command line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If host begins with a
slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the Internet TCP/IP port or local Unix domain socket file extension on which the server
is listening for connections.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password
is not available by other means such as a .pgpass file, the connection attempt will fail. This
option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force droplang to prompt for a password before connecting to a database.

This option is never essential, since droplang will automatically prompt for a password if the
server demands password authentication. However, droplang will waste a connection attempt
finding out that the server wants a password. In some cases it is worth typing -W to avoid the
extra connection attempt.

Environment

PGDATABASE
PGHOST
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by
libpq (see Section 32.14).
Diagnostics
Most error messages are self-explanatory. If not, run droplang with the \texttt{--echo} option and see under the respective SQL command for details. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Notes
Use createlang to add a language.

Examples
To remove the language pltcl:

\$ droplang pltcl dbname

See Also
createlang, DROP EXTENSION, DROP LANGUAGE
dropuser

Name

dropuser — remove a PostgreSQL user account

Synopsis

dropuser [connection-option...][option...][username]

Description

dropuser removes an existing PostgreSQL user. Only superusers and users with the CREATEROLE privilege can remove PostgreSQL users. (To remove a superuser, you must yourself be a superuser.) dropuser is a wrapper around the SQL command DROP ROLE. There is no effective difference between dropping users via this utility and via other methods for accessing the server.

Options

dropuser accepts the following command-line arguments:

username

   Specifies the name of the PostgreSQL user to be removed. You will be prompted for a name if none is specified on the command line and the -i/--interactive option is used.
-e
--echo

   Echo the commands that dropuser generates and sends to the server.
-i
--interactive

   Prompt for confirmation before actually removing the user, and prompt for the user name if none is specified on the command line.
-V
--version

   Print the dropuser version and exit.
--if-exists

   Do not throw an error if the user does not exist. A notice is issued in this case.
-?
--help

   Show help about dropuser command line arguments, and exit.

dropuser also accepts the following command-line arguments for connection parameters:
-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a
slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening
for connections.

-U username
--username=username

User name to connect as (not the user name to drop).

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password
is not available by other means such as a .pgpass file, the connection attempt will fail. This
option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force dropuser to prompt for a password before connecting to a database.

This option is never essential, since dropuser will automatically prompt for a password if the
server demands password authentication. However, dropuser will waste a connection attempt
finding out that the server wants a password. In some cases it is worth typing -W to avoid the
extra connection attempt.

Environment

PGHOST
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by
libpq (see Section 32.14).

Diagnostics

In case of difficulty, see DROP ROLE and psql for discussions of potential problems and error mes-
riages. The database server must be running at the targeted host. Also, any default connection settings
and environment variables used by the libpq front-end library will apply.
**Examples**

To remove user *joe* from the default database server:

```bash
$ dropuser joe
```

To remove user *joe* using the server on host *eden*, port 5000, with verification and a peek at the underlying command:

```bash
$ dropuser -p 5000 -h eden -i -e joe
Role "joe" will be permanently removed.
Are you sure? (y/n) y
DROP ROLE joe;
```

**See Also**

`createuser`, `DROP ROLE`
ecpg

Name

ecpg — embedded SQL C preprocessor

Synopsis

ecpg [option...] file...

Description

ecpg is the embedded SQL preprocessor for C programs. It converts C programs with embedded SQL statements to normal C code by replacing the SQL invocations with special function calls. The output files can then be processed with any C compiler tool chain.

ecpg will convert each input file given on the command line to the corresponding C output file. Input files preferably have the extension .pgc. The extension will be replaced by .c to determine the output file name. The output file name can also be overridden using the -o option.

This reference page does not describe the embedded SQL language. See Chapter 34 for more information on that topic.

Options

ecpg accepts the following command-line arguments:

- c
  Automatically generate certain C code from SQL code. Currently, this works for EXEC SQL TYPE.
- C mode
  Set a compatibility mode. mode can be INFORMIX or INFORMIX_SE.
- D symbol
  Define a C preprocessor symbol.
- I directory
  Parse system include files as well.
- I directory
  Specify an additional include path, used to find files included via EXEC SQL INCLUDE. Defaults are . (current directory), /usr/local/include, the PostgreSQL include directory which is defined at compile time (default: /usr/local/pgsql/include), and /usr/include, in that order.
- o filename
  Specifies that ecpg should write all its output to the given filename.
-r option

Selects run-time behavior. Option can be one of the following:

no_indicator

Do not use indicators but instead use special values to represent null values. Historically there have been databases using this approach.

prepare

Prepare all statements before using them. Libecpg will keep a cache of prepared statements and reuse a statement if it gets executed again. If the cache runs full, libecpg will free the least used statement.

questionmarks

Allow question mark as placeholder for compatibility reasons. This used to be the default long ago.

-t

Turn on autocommit of transactions. In this mode, each SQL command is automatically committed unless it is inside an explicit transaction block. In the default mode, commands are committed only when EXEC SQL COMMIT is issued.

-v

Print additional information including the version and the "include" path.

--version

Print the ecpg version and exit.

-?

--help

Show help about ecpg command line arguments, and exit.

Notes

When compiling the preprocessed C code files, the compiler needs to be able to find the ECPG header files in the PostgreSQL include directory. Therefore, you might have to use the -I option when invoking the compiler (e.g., -I/usr/local/pgsql/include).

Programs using C code with embedded SQL have to be linked against the libecpg library, for example using the linker options -L/usr/local/pgsql/lib -lecpg.

The value of either of these directories that is appropriate for the installation can be found out using pg_config.

Examples

If you have an embedded SQL C source file named prog1.pgc, you can create an executable program using the following sequence of commands:

ecpg prog1.pgc

1791
cc -I/usr/local/pgsql/include -c progl.c
cc -o progl progl.o -L/usr/local/pgsql/lib -lcp

ecpg
**pg_basebackup**

**Name**

`pg_basebackup` — take a base backup of a PostgreSQL cluster

**Synopsis**

`pg_basebackup [option...]`

**Description**

`pg_basebackup` is used to take base backups of a running PostgreSQL database cluster. These are taken without affecting other clients to the database, and can be used both for point-in-time recovery (see Section 25.3) and as the starting point for a log shipping or streaming replication standby servers (see Section 26.2).

`pg_basebackup` makes a binary copy of the database cluster files, while making sure the system is put in and out of backup mode automatically. Backups are always taken of the entire database cluster; it is not possible to back up individual databases or database objects. For individual database backups, a tool such as `pg_dump` must be used.

The backup is made over a regular PostgreSQL connection, and uses the replication protocol. The connection must be made with a superuser or a user having `REPLICATION` permissions (see Section 21.2), and `pg_hba.conf` must explicitly permit the replication connection. The server must also be configured with `max_wal_senders` set high enough to leave at least one session available for the backup.

There can be multiple `pg_basebackup`s running at the same time, but it is better from a performance point of view to take only one backup, and copy the result.

`pg_basebackup` can make a base backup from not only the master but also the standby. To take a backup from the standby, set up the standby so that it can accept replication connections (that is, set `max_wal_senders` and `hot_standby`, and configure host-based authentication). You will also need to enable `full_page_writes` on the master.

Note that there are some limitations in an online backup from the standby:

- The backup history file is not created in the database cluster backed up.
- There is no guarantee that all WAL files required for the backup are archived at the end of backup. If you are planning to use the backup for an archive recovery and want to ensure that all required files are available at that moment, you need to include them into the backup by using the `-x` option.
- If the standby is promoted to the master during online backup, the backup fails.
- All WAL records required for the backup must contain sufficient full-page writes, which requires you to enable `full_page_writes` on the master and not to use a tool like `pg_compresslog` as `archive_command` to remove full-page writes from WAL files.
Options

The following command-line options control the location and format of the output.

-D directory
--pgdata=directory

Directory to write the output to. pg_basebackup will create the directory and any parent directories if necessary. The directory may already exist, but it is an error if the directory already exists and is not empty.

When the backup is in tar mode, and the directory is specified as -(dash), the tar file will be written to stdout.

This option is required.

-F format
--format=format

Selects the format for the output. format can be one of the following:

p
plain

Write the output as plain files, with the same layout as the current data directory and tablespaces. When the cluster has no additional tablespaces, the whole database will be placed in the target directory. If the cluster contains additional tablespaces, the main data directory will be placed in the target directory, but all other tablespaces will be placed in the same absolute path as they have on the server.

This is the default format.

t

tar

Write the output as tar files in the target directory. The main data directory will be written to a file named base.tar, and all other tablespaces will be named after the tablespace OID.

If the value -(dash) is specified as target directory, the tar contents will be written to standard output, suitable for piping to for example gzip. This is only possible if the cluster has no additional tablespaces.

-r rate
--max-rate=rate

The maximum transfer rate of data transferred from the server. Values are in kilobytes per second. Use a suffix of M to indicate megabytes per second. A suffix of k is also accepted, and has no effect. Valid values are between 32 kilobytes per second and 1024 megabytes per second.

The purpose is to limit the impact of pg_basebackup on the running server.

This option always affects transfer of the data directory. Transfer of WAL files is only affected if the collection method is fetch.

-R
--write-recovery-conf

Write a minimal recovery.conf in the output directory (or into the base archive file when using tar format) to ease setting up a standby server. The recovery.conf file will record the connection settings and, if specified, the replication slot that pg_basebackup is using, so that the streaming replication will use the same settings later on.


```
-S slotname
--slot=slotname

This option can only be used together with -X stream. It causes the WAL streaming to use
the specified replication slot. If the base backup is intended to be used as a streaming replication
standby using replication slots, it should then use the same replication slot name in
recovery.conf. That way, it is ensured that the server does not remove any necessary WAL
data in the time between the end of the base backup and the start of streaming replication.

-T olddir=newdir
--tablespace-mapping=olddir=newdir

Relocate the tablespace in directory olddir to newdir during the backup. To be effective, 
olddir must exactly match the path specification of the tablespace as it is currently defined. 
(But it is not an error if there is no tablespace in olddir contained in the backup.) Both olddir 
and newdir must be absolute paths. If a path happens to contain a = sign, escape it with a \backslash. This option can be specified multiple times for multiple tablespaces. See examples 
below.

If a tablespace is relocated in this way, the symbolic links inside the main data directory are 
updated to point to the new location. So the new data directory is ready to be used for a new 
server instance with all tablespaces in the updated locations.

--xlogdir=xlogdir

Specifies the location for the transaction log directory. xlogdir must be an absolute path. The 
transaction log directory can only be specified when the backup is in plain mode.

-x
--xlog

Using this option is equivalent of using -X with method fetch.

-X method
--xlog-method=method

Includes the required transaction log files (WAL files) in the backup. This will include all transac-
tion logs generated during the backup. If this option is specified, it is possible to start a postmaster 
directly in the extracted directory without the need to consult the log archive, thus making this a 
completely standalone backup.

The following methods for collecting the transaction logs are supported:

f
fetch

The transaction log files are collected at the end of the backup. Therefore, it is necessary for 
the wal_keep_segments parameter to be set high enough that the log is not removed before 
the end of the backup. If the log has been rotated when it’s time to transfer it, the backup 
will fail and be unusable.

s
stream

Stream the transaction log while the backup is created. This will open a second connection 
to the server and start streaming the transaction log in parallel while running the backup. 
Therefore, it will use up two connections configured by the max_wal_senders parameter. 
As long as the client can keep up with transaction log received, using this mode requires no 
extra transaction logs to be saved on the master.
```
The following command-line options control the generation of the backup and the running of the program.

```
-z
--gzip
```

Enables gzip compression of tar file output, with the default compression level. Compression is only available when using the tar format.

```
-Z level
--compress=level
```

Enables gzip compression of tar file output, and specifies the compression level (0 through 9, 0 being no compression and 9 being best compression). Compression is only available when using the tar format.

The following command-line options control the database connection parameters.

```
-d connstr
--dbname=connstr
```

Specifies parameters used to connect to the server, as a connection string. See Section 32.1.1 for more information.

The option is called `--dbname` for consistency with other client applications, but because `pg_basebackup` doesn’t connect to any particular database in the cluster, database name in the connection string will be ignored.
pg_basebackup

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

-s interval
--status-interval=interval

Specifies the number of seconds between status packets sent back to the server. This allows for easier monitoring of the progress from server. A value of zero disables the periodic status updates completely, although an update will still be sent when requested by the server, to avoid timeout disconnect. The default value is 10 seconds.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force pg_basebackup to prompt for a password before connecting to a database.

This option is never essential, since pg_basebackup will automatically prompt for a password if the server demands password authentication. However, pg_basebackup will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

Other options are also available:

-V
--version

Print the pg_basebackup version and exit.

-?
--help

Show help about pg_basebackup command line arguments, and exit.
Environment

This utility, like most other PostgreSQL utilities, uses the environment variables supported by libpq (see Section 32.14).

Notes

At the beginning of the backup, a checkpoint needs to be written on the server the backup is taken from. Especially if the option --checkpoint-fast is not used, this can take some time during which pg_basebackup will be appear to be idle.

The backup will include all files in the data directory and tablespaces, including the configuration files and any additional files placed in the directory by third parties. But only regular files and directories are copied. Symbolic links (other than those used for tablespaces) and special device files are skipped. (See Section 51.3 for the precise details.)

Tablespaces will in plain format by default be backed up to the same path they have on the server, unless the option --tablespace-mapping is used. Without this option, running a plain format base backup on the same host as the server will not work if tablespaces are in use, because the backup would have to be written to the same directory locations as the original tablespaces.

When tar format mode is used, it is the user’s responsibility to unpack each tar file before starting the PostgreSQL server. If there are additional tablespaces, the tar files for them need to be unpacked in the correct locations. In this case the symbolic links for those tablespaces will be created by the server according to the contents of the tablespace_map file that is included in the base.tar file.

pg_basebackup works with servers of the same or an older major version, down to 9.1. However, WAL streaming mode (-X stream) only works with server version 9.3 and later, and tar format mode (--format=tar) of the current version only works with server version 9.5 or later.

Examples

To create a base backup of the server at mydbserver and store it in the local directory /usr/local/pgsql/data:

```
$ pg_basebackup -h mydbserver -D /usr/local/pgsql/data
```

To create a backup of the local server with one compressed tar file for each tablespace, and store it in the directory backup, showing a progress report while running:

```
$ pg_basebackup -D backup -Ft -z -P
```

To create a backup of a single-tablespace local database and compress this with bzip2:

```
$ pg_basebackup -D - -Ft | bzip2 > backup.tar.bz2
```

(This command will fail if there are multiple tablespaces in the database.)

To create a backup of a local database where the tablespace in /opt/ts is relocated to ./backup/ts:

```
$ pg_basebackup -D backup/data -T /opt/ts=$(pwd)/backup/ts
```

1798
See Also

pg_dump
pgbench

Name
pgbench — run a benchmark test on PostgreSQL

Synopsis
pgbench -i [option...] [dbname]

pgbench [option...] [dbname]

Description
pgbench is a simple program for running benchmark tests on PostgreSQL. It runs the same sequence of SQL commands over and over, possibly in multiple concurrent database sessions, and then calculates the average transaction rate (transactions per second). By default, pgbench tests a scenario that is loosely based on TPC-B, involving five SELECT, UPDATE, and INSERT commands per transaction. However, it is easy to test other cases by writing your own transaction script files.

Typical output from pgbench looks like:

transaction type: <builtin: TPC-B (sort of)>
scaling factor: 10
query mode: simple
number of clients: 10
number of threads: 1
number of transactions per client: 1000
number of transactions actually processed: 10000/10000
tps = 85.184871 (including connections establishing)
tps = 85.296346 (excluding connections establishing)

The first six lines report some of the most important parameter settings. The next line reports the number of transactions completed and intended (the latter being just the product of number of clients and number of transactions per client); these will be equal unless the run failed before completion. (In -T mode, only the actual number of transactions is printed.) The last two lines report the number of transactions per second, figured with and without counting the time to start database sessions.

The default TPC-B-like transaction test requires specific tables to be set up beforehand. pgbench should be invoked with the -i (initialize) option to create and populate these tables. (When you are testing a custom script, you don’t need this step, but will instead need to do whatever setup your test needs.) Initialization looks like:

pgbench -i [ other-options ] dbname

where dbname is the name of the already-created database to test in. (You may also need -h, -p, and/or -U options to specify how to connect to the database server.)
Caution

pgbench -i creates four tables pgbench_accounts, pgbench_branches, pgbench_history, and pgbench_tellers, destroying any existing tables of these names. Be very careful to use another database if you have tables having these names!

At the default “scale factor” of 1, the tables initially contain this many rows:

<table>
<thead>
<tr>
<th>table</th>
<th># of rows</th>
</tr>
</thead>
<tbody>
<tr>
<td>pgbench_branches</td>
<td>1</td>
</tr>
<tr>
<td>pgbench_tellers</td>
<td>10</td>
</tr>
<tr>
<td>pgbench_accounts</td>
<td>100000</td>
</tr>
<tr>
<td>pgbench_history</td>
<td>0</td>
</tr>
</tbody>
</table>

You can (and, for most purposes, probably should) increase the number of rows by using the -s (scale factor) option. The -F (fillfactor) option might also be used at this point.

Once you have done the necessary setup, you can run your benchmark with a command that doesn’t include -i, that is

```
pgbench [ options ] dbname
```

In nearly all cases, you’ll need some options to make a useful test. The most important options are -c (number of clients), -t (number of transactions), -T (time limit), and -f (specify a custom script file). See below for a full list.

**Options**

The following is divided into three subsections: Different options are used during database initialization and while running benchmarks, some options are useful in both cases.

**Initialization Options**

pgbench accepts the following command-line initialization arguments:

- `-i`
  --initialize
  Required to invoke initialization mode.

- `-F fillfactor`
  --fillfactor=fillfactor
  Create the pgbench_accounts, pgbench_tellers and pgbench_branches tables with the given fillfactor. Default is 100.

- `-n`
  --no-vacuum
  Perform no vacuuming after initialization.
-q
--quiet
Switch logging to quiet mode, producing only one progress message per 5 seconds. The default logging prints one message each 100000 rows, which often outputs many lines per second (especially on good hardware).

-s scale_factor
--scale=scale_factor
Multiply the number of rows generated by the scale factor. For example, -s 100 will create 10,000,000 rows in the pgbench_accounts table. Default is 1. When the scale is 20,000 or larger, the columns used to hold account identifiers (aid columns) will switch to using larger integers (bigint), in order to be big enough to hold the range of account identifiers.

--foreign-keys
Create foreign key constraints between the standard tables.

--index-tablespace=index_tablespace
Create indexes in the specified tablespace, rather than the default tablespace.

--tablespace=tablespace
Create tables in the specified tablespace, rather than the default tablespace.

--unlogged-tables
Create all tables as unlogged tables, rather than permanent tables.

**Benchmarking Options**

pgbench accepts the following command-line benchmarking arguments:

-b scriptname[@weight]
--builtin=scriptname[@weight]
Add the specified built-in script to the list of executed scripts. An optional integer weight after @ allows to adjust the probability of drawing the script. If not specified, it is set to 1. Available built-in scripts are: tpcb-like, simple-update and select-only. Unambiguous prefixes of built-in names are accepted. With special name list, show the list of built-in scripts and exit immediately.

-c clients
--client=clients
Number of clients simulated, that is, number of concurrent database sessions. Default is 1.

-C
--connect
Establish a new connection for each transaction, rather than doing it just once per client session. This is useful to measure the connection overhead.

-d
--debug
Print debugging output.
pgbench

-D varname=value
--define=varname=value

Define a variable for use by a custom script (see below). Multiple -D options are allowed.

-f filename[@weight]
--file=filename[@weight]

Add a transaction script read from filename to the list of executed scripts. An optional integer weight after @ allows to adjust the probability of drawing the test. See below for details.

-j threads
--jobs=threads

Number of worker threads within pgbench. Using more than one thread can be helpful on multi-CPU machines. Clients are distributed as evenly as possible among available threads. Default is 1.

-l
--log

Write the time taken by each transaction to a log file. See below for details.

-L limit
--latency-limit=limit

Transactions that last more than limit milliseconds are counted and reported separately, as late. When throttling is used (--rate=...), transactions that lag behind schedule by more than limit ms, and thus have no hope of meeting the latency limit, are not sent to the server at all. They are counted and reported separately as skipped.

-M querymode
--protocol=querymode

Protocol to use for submitting queries to the server:

• simple: use simple query protocol.
• extended: use extended query protocol.
• prepared: use extended query protocol with prepared statements.

The default is simple query protocol. (See Chapter 51 for more information.)

-n
--no-vacuum

Perform no vacuuming before running the test. This option is necessary if you are running a custom test scenario that does not include the standard tables pgbench_accounts, pgbench_branches, pgbench_history, and pgbench_tellers.

-N
--skip-some-updates


-P sec
--progress=sec

Show progress report every sec seconds. The report includes the time since the beginning of the run, the tps since the last report, and the transaction latency average and standard deviation since the last report. Under throttling (--R), the latency is computed with respect to the transaction scheduled start time, not the actual transaction beginning time, thus it also includes the average schedule lag time.
pgbench

-\( r \)
--report-latencies

Report the average per-statement latency (execution time from the perspective of the client) of each command after the benchmark finishes. See below for details.

-\( R \) rate
--rate=rate

Execute transactions targeting the specified rate instead of running as fast as possible (the default). The rate is given in transactions per second. If the targeted rate is above the maximum possible rate, the rate limit won’t impact the results.

The rate is targeted by starting transactions along a Poisson-distributed schedule time line. The expected start time schedule moves forward based on when the client first started, not when the previous transaction ended. That approach means that when transactions go past their original scheduled end time, it is possible for later ones to catch up again.

When throttling is active, the transaction latency reported at the end of the run is calculated from the scheduled start times, so it includes the time each transaction had to wait for the previous transaction to finish. The wait time is called the schedule lag time, and its average and maximum are also reported separately. The transaction latency with respect to the actual transaction start time, i.e. the time spent executing the transaction in the database, can be computed by subtracting the schedule lag time from the reported latency.

If --latency-limit is used together with --rate, a transaction can lag behind so much that it is already over the latency limit when the previous transaction ends, because the latency is calculated from the scheduled start time. Such transactions are not sent to the server, but are skipped altogether and counted separately.

A high schedule lag time is an indication that the system cannot process transactions at the specified rate, with the chosen number of clients and threads. When the average transaction execution time is longer than the scheduled interval between each transaction, each successive transaction will fall further behind, and the schedule lag time will keep increasing the longer the test run is. When that happens, you will have to reduce the specified transaction rate.

-\( s \) scale_factor
--scale=scale_factor

Report the specified scale factor in pgbench’s output. With the built-in tests, this is not necessary; the correct scale factor will be detected by counting the number of rows in the pgbench_branches table. However, when testing only custom benchmarks (-f option), the scale factor will be reported as 1 unless this option is used.

-S
--select-only

Run built-in select-only script. Shorthand for -b select-only.

-\( t \) transactions
--transactions=transactions

Number of transactions each client runs. Default is 10.

-\( T \) seconds
--time=seconds

Run the test for this many seconds, rather than a fixed number of transactions per client. -t and -T are mutually exclusive.
Vacuum all four standard tables before running the test. With neither \texttt{-n} nor \texttt{-v}, \texttt{pgbench} will vacuum the \texttt{pgbench_tellers} and \texttt{pgbench_branches} tables, and will truncate \texttt{pgbench_history}.

\textbf{--aggregate-interval=seconds}

Length of aggregation interval (in seconds). May be used only together with \texttt{-l} - with this option, the log contains per-interval summary (number of transactions, min/max latency and two additional fields useful for variance estimation).

This option is not currently supported on Windows.

\textbf{--progress-timestamp}

When showing progress (option \texttt{-P}), use a timestamp (Unix epoch) instead of the number of seconds since the beginning of the run. The unit is in seconds, with millisecond precision after the dot. This helps compare logs generated by various tools.

\textbf{--sampling-rate=rate}

Sampling rate, used when writing data into the log, to reduce the amount of log generated. If this option is given, only the specified fraction of transactions are logged. 1.0 means all transactions will be logged, 0.05 means only 5\% of the transactions will be logged.

Remember to take the sampling rate into account when processing the log file. For example, when computing \texttt{tps} values, you need to multiply the numbers accordingly (e.g. with 0.01 sample rate, you’ll only get 1/100 of the actual \texttt{tps}).

\section*{Common Options}

\texttt{pgbench} accepts the following command-line common arguments:

\texttt{-h hostname}
\texttt{--host=hostname}

The database server’s host name

\texttt{-p port}
\texttt{--port=port}

The database server’s port number

\texttt{-U login}
\texttt{--username=login}

The user name to connect as

\texttt{-v}
\texttt{--version}

Print the \texttt{pgbench} version and exit.

\texttt{-?}
\texttt{--help}

Show help about \texttt{pgbench} command line arguments, and exit.
Notes

What is the “Transaction” Actually Performed in pgbench?

pgbench executes test scripts chosen randomly from a specified list. They include built-in scripts with -b and user-provided custom scripts with -f. Each script may be given a relative weight specified after a @ so as to change its drawing probability. The default weight is 1. Scripts with a weight of 0 are ignored.

The default built-in transaction script (also invoked with -b tpcb-like) issues seven commands per transaction over randomly chosen aid, tid, bid and delta. The scenario is inspired by the TPC-B benchmark, but is not actually TPC-B, hence the name.

1. BEGIN;
2. UPDATE pgbench_accounts SET abalance = abalance + :delta WHERE aid = :aid;
3. SELECT abalance FROM pgbench_accounts WHERE aid = :aid;
4. UPDATE pgbench_tellers SET tbalance = tbalance + :delta WHERE tid = :tid;
5. UPDATE pgbench_branches SET bbalance = bbalance + :delta WHERE bid = :bid;
6. INSERT INTO pgbench_history (tid, bid, aid, delta, mtime) VALUES (:tid, :bid, :aid, :delta, CURRENT_TIMESTAMP);
7. END;

If you select the simple-update built-in (also -N), steps 4 and 5 aren’t included in the transaction. This will avoid update contention on these tables, but it makes the test case even less like TPC-B.

If you select the select-only built-in (also -S), only the SELECT is issued.

Custom Scripts

pgbench has support for running custom benchmark scenarios by replacing the default transaction script (described above) with a transaction script read from a file (-f option). In this case a “transaction” counts as one execution of a script file.

A script file contains one or more SQL commands terminated by semicolons. Empty lines and lines beginning with -- are ignored. Script files can also contain “meta commands”, which are interpreted by pgbench itself, as described below.

Note: Before PostgreSQL 9.6, SQL commands in script files were terminated by newlines, and so they could not be continued across lines. Now a semicolon is required to separate consecutive SQL commands (though a SQL command does not need one if it is followed by a meta command).

If you need to create a script file that works with both old and new versions of pgbench, be sure to write each SQL command on a single line ending with a semicolon.

There is a simple variable-substitution facility for script files. Variables can be set by the command-line -D option, explained above, or by the meta commands explained below. In addition to any variables preset by -D command-line options, there are a few variables that are preset automatically, listed in Table 1. A value specified for these variables using -D takes precedence over the automatic presets.
Once set, a variable’s value can be inserted into a SQL command by writing `:variablename`. When running more than one client session, each session has its own set of variables.

### Table 1. Automatic Variables

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>scale</td>
<td>current scale factor</td>
</tr>
<tr>
<td>client_id</td>
<td>unique number identifying the client session</td>
</tr>
<tr>
<td></td>
<td>(starts from zero)</td>
</tr>
</tbody>
</table>

Script file meta commands begin with a backslash (\) and extend to the end of the line. Arguments to a meta command are separated by white space. These meta commands are supported:

\set varname expression

Sets variable `varname` to a value calculated from `expression`. The expression may contain integer constants such as 5432, double constants such as 3.14159, references to variables `:variablename`, unary operators (+, −) and binary operators (+, −, *, /, %) with their usual precedence and associativity, function calls, and parentheses.

Examples:

\set ntellers 10 * :scale  
\set aid (1021 * random(1, 100000 * :scale)) % (100000 * :scale) + 1

\sleep number [ us | ms | s ]

Causes script execution to sleep for the specified duration in microseconds (us), milliseconds (ms) or seconds (s). If the unit is omitted then seconds are the default. `number` can be either an integer constant or a `:variablename` reference to a variable having an integer value.

Example:

\sleep 10 ms

\setshell varname command [ argument ... ]

Sets variable `varname` to the result of the shell command `command` with the given `argument(s)`. The command must return an integer value through its standard output.

`command` and each `argument` can be either a text constant or a `:variablename` reference to a variable. If you want to use an `argument` starting with a colon, write an additional colon at the beginning of `argument`.

Example:

\setshell variable_to_be_assigned command literal_argument :variable ::literal_starting_with_colon

\shell command [ argument ... ]

Same as `\setshell`, but the result of the command is discarded.

Example:

\shell command literal_argument :variable ::literal_starting_with_colon

### Built-In Functions

The functions listed in Table 2 are built into pgbench and may be used in expressions appearing in `\set`.

1807
Table 2. pgbench Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>abs(a)</td>
<td>same as a</td>
<td>absolute value</td>
<td>abs(-17)</td>
<td>17</td>
</tr>
<tr>
<td>debug(a)</td>
<td>same as a</td>
<td>print a to stderr, and return a</td>
<td>debug(5432.1)</td>
<td>5432.1</td>
</tr>
<tr>
<td>double(i)</td>
<td>double</td>
<td>cast to double</td>
<td>double(5432)</td>
<td>5432.0</td>
</tr>
<tr>
<td>greatest(a [, ... ] )</td>
<td>double if any a is double, else integer</td>
<td>largest value among arguments</td>
<td>greatest(5, 4, 3, 2)</td>
<td>5</td>
</tr>
<tr>
<td>int(x)</td>
<td>integer</td>
<td>cast to int</td>
<td>int(5.4 + 3.8)</td>
<td>9</td>
</tr>
<tr>
<td>least(a [, ... ] )</td>
<td>double if any a is double, else integer</td>
<td>smallest value among arguments</td>
<td>least(5, 4, 3, 2.1)</td>
<td>2.1</td>
</tr>
<tr>
<td>pi()</td>
<td>double</td>
<td>value of the constant PI</td>
<td>pi()</td>
<td>3.14159265358979323846</td>
</tr>
<tr>
<td>random(lb, ub)</td>
<td>integer</td>
<td>uniformly-distributed random integer in [lb, ub]</td>
<td>random(1, 10)</td>
<td>an integer between 1 and 10</td>
</tr>
<tr>
<td>random_exponential(lb, ub, parameter)</td>
<td>integer</td>
<td>exponentially-distributed random integer in [lb, ub], see below</td>
<td>random_exponential(10, 3.0)</td>
<td>an integer between 1 and 10</td>
</tr>
<tr>
<td>random_gaussian(lb, ub, parameter)</td>
<td>integer</td>
<td>Gaussian-distributed random integer in [lb, ub], see below</td>
<td>random_gaussian(10, 2.5)</td>
<td>an integer between 1 and 10</td>
</tr>
<tr>
<td>sqrt(x)</td>
<td>double</td>
<td>square root</td>
<td>sqrt(2.0)</td>
<td>1.414213562</td>
</tr>
</tbody>
</table>

The random function generates values using a uniform distribution, that is all the values are drawn within the specified range with equal probability. The random_exponential and random_gaussian functions require an additional double parameter which determines the precise shape of the distribution.

- For an exponential distribution, parameter controls the distribution by truncating a quickly-decreasing exponential distribution at parameter, and then projecting onto integers between the bounds. To be precise, with

  \[ f(x) = \text{exp}(-\text{parameter} \cdot (x - \text{min}) / (\text{max} - \text{min} + 1)) / (1 - \text{exp}(-\text{parameter})) \]

  Then value \( i \) between \( \text{min} \) and \( \text{max} \) inclusive is drawn with probability: \( f(i) - f(i + 1) \).

  Intuitively, the larger the parameter, the more frequently values close to \( \text{min} \) are accessed, and the less frequently values close to \( \text{max} \) are accessed. The closer to 0 parameter is, the flatter (more uniform) the access distribution. A crude approximation of the distribution is that the most frequent 1% values in the range, close to \( \text{min} \), are drawn parameter% of the time. The parameter value must be strictly positive.
For a Gaussian distribution, the interval is mapped onto a standard normal distribution (the classical bell-shaped Gaussian curve) truncated at \(-\text{parameter}\) on the left and \(+\text{parameter}\) on the right. Values in the middle of the interval are more likely to be drawn. To be precise, if \(\Phi(x)\) is the cumulative distribution function of the standard normal distribution, with mean \(\mu\) defined as \((\max + \min) / 2.0\), with

\[
f(x) = \Phi(2.0 * \text{parameter} * (x - \mu) / (\max - \min + 1)) / (2.0 * \Phi(\text{parameter}) - 1)
\]

then value \(i\) between \(\min\) and \(\max\) inclusive is drawn with probability: \(f(i + 0.5) - f(i - 0.5)\). Intuitively, the larger the \(\text{parameter}\), the more frequently values close to the middle of the interval are drawn, and the less frequently values close to the \(\min\) and \(\max\) bounds. About 67% of values are drawn from the middle 1.0 / \(\text{parameter}\), that is a relative 0.5 / \(\text{parameter}\) around the mean, and 95% in the middle 2.0 / \(\text{parameter}\), that is a relative 1.0 / \(\text{parameter}\) around the mean; for instance, if \(\text{parameter}\) is 4.0, 67% of values are drawn from the middle quarter (1.0 / 4.0) of the interval (i.e. from 3.0 / 8.0 to 5.0 / 8.0) and 95% from the middle half (2.0 / 4.0) of the interval (second and third quartiles). The minimum \(\text{parameter}\) is 2.0 for performance of the Box-Muller transform.

As an example, the full definition of the built-in TPC-B-like transaction is:

```sql
\set aid random(1, 100000 * :scale)
\set bid random(1, 1 * :scale)
\set tid random(1, 10 * :scale)
\set delta random(-5000, 5000)
BEGIN;
  UPDATE pgbench_accounts SET abalance = abalance + :delta WHERE aid = :aid;
  SELECT abalance FROM pgbench_accounts WHERE aid = :aid;
  UPDATE pgbench_tellers SET tbalance = tbalance + :delta WHERE tid = :tid;
  UPDATE pgbench_branches SET bbalance = bbalance + :delta WHERE bid = :bid;
  INSERT INTO pgbench_history (tid, bid, aid, delta, mtime) VALUES (:tid, :bid, :aid, :delta, CURRENT_TIMESTAMP);
END;
```

This script allows each iteration of the transaction to reference different, randomly-chosen rows. (This example also shows why it’s important for each client session to have its own variables — otherwise they’d not be independently touching different rows.)

### Per-Transaction Logging

With the \(-l\) option but without the \(--aggregate-interval\), \texttt{pgbench} writes the time taken by each transaction to a log file. The log file will be named \texttt{pgbench_log.nnn}, where \(nnn\) is the PID of the \texttt{pgbench} process. If the \(-j\) option is 2 or higher, creating multiple worker threads, each will have its own log file. The first worker will use the same name for its log file as in the standard single worker case. The additional log files for the other workers will be named \texttt{pgbench_log.nnn.mmm}, where \(mmm\) is a sequential number for each worker starting with 1.

The format of the log is:

```
client_id transaction_no time script_no time_epoch time_us [schedule_lag]
```

where \(time\) is the total elapsed transaction time in microseconds, \(script_no\) identifies which script file was used (useful when multiple scripts were specified with \(-f\) or \(-b\)), and \(time_epoch\) and \(time_us\) are a Unix epoch format time stamp and an offset in microseconds (suitable for creating an ISO 8601 time stamp with fractional seconds) showing when the transaction completed. Field \(schedule_lag\) is the difference between the transaction’s scheduled start time, and the time it actually started, in microseconds. It is only present when the \(--rate\) option is used.
When both `--rate` and `--latency-limit` are used, the time for a skipped transaction will be reported as skipped.

Here is a snippet of the log file generated:

```
0 199 2241 0 1175850568 995598
0 200 2465 0 1175850568 998079
0 201 2513 0 1175850569 608
0 202 2038 0 1175850569 2663
```

Another example with `--rate=100` and `--latency-limit=5` (note the additional `schedule_lag` column):

```
0 81 4621 0 1412881037 912698 3005
0 82 6173 0 1412881037 914578 4304
0 83 skipped 0 1412881037 914578 5217
0 83 skipped 0 1412881037 914578 5099
0 83 4722 0 1412881037 916203 3108
0 84 4142 0 1412881037 918023 2333
0 85 2465 0 1412881037 919759 740
```

In this example, transaction 82 was late, because its latency (6.173 ms) was over the 5 ms limit. The next two transactions were skipped, because they were already late before they were even started.

When running a long test on hardware that can handle a lot of transactions, the log files can become very large. The `--sampling-rate` option can be used to log only a random sample of transactions.

### Aggregated Logging

With the `--aggregate-interval` option, the logs use a bit different format:

```
interval_start num_of_transactions latency_sum latency_2_sum min_latency max_latency [lag_sum lag_2_sum min_lag max_lag] skipped_transactions
```

where `interval_start` is the start of the interval (Unix epoch format time stamp), `num_of_transactions` is the number of transactions within the interval, `latency_sum` is a sum of latencies (so you can compute average latency easily). The following two fields are useful for variance estimation - `latency_sum` is a sum of latencies and `latency_2_sum` is a sum of 2nd powers of latencies. The next two fields are `min_latency` - a minimum latency within the interval, and `max_latency` - maximum latency within the interval. A transaction is counted into the interval when it was committed. The fields in the end, `lag_sum`, `lag_2_sum`, `min_lag`, and `max_lag`, are only present if the `--rate` option is used. The very last one, `skipped_transactions`, is only present if the option `--latency-limit` is present, too. They are calculated from the time each transaction had to wait for the previous one to finish, i.e. the difference between each transaction’s scheduled start time and the time it actually started.

Here is example output:

```
1345828501 5601 1542744 483552416 61 2573
1345828503 7884 1979812 565806736 60 1479
1345828505 7208 1979422 567277552 59 1391
1345828507 7685 1980268 569784714 60 1398
1345828509 7073 1979779 573489941 236 1411
```

Notice that while the plain (unaggregated) log file contains a reference to the custom script files, the aggregated log does not. Therefore if you need per script data, you need to aggregate the data on your own.
Per-Statement Latencies

With the `-r` option, pgbench collects the elapsed transaction time of each statement executed by every client. It then reports an average of those values, referred to as the latency for each statement, after the benchmark has finished.

For the default script, the output will look similar to this:

```
starting vacuum...end.
transaction type: <builtin: TPC-B (sort of)>
scaling factor: 1
query mode: simple
number of clients: 10
number of threads: 1
number of transactions per client: 1000
number of transactions actually processed: 10000/10000
latency average = 15.844 ms
latency stddev = 2.715 ms
tps = 618.764555 (including connections establishing)
tps = 622.977698 (excluding connections establishing)
script statistics:
- statement latencies in milliseconds:
  0.002 \set aid random(1, 100000 * :scale)
  0.005 \set bid random(1, 1 * :scale)
  0.002 \set tid random(1, 10 * :scale)
  0.001 \set delta random(-5000, 5000)
  0.326 BEGIN;
  0.603 UPDATE pgbench_accounts SET abalance = abalance + :delta WHERE aid = :aid;
  0.454 SELECT abalance FROM pgbench_accounts WHERE aid = :aid;
  5.528 UPDATE pgbench_tellers SET tbalance = tbalance + :delta WHERE tid = :tid;
  7.335 UPDATE pgbench_branches SET bbalance = bbalance + :delta WHERE bid = :bid;
  0.371 INSERT INTO pgbench_history (tid, bid, aid, delta, mtime) VALUES (:tid, :bid, :aid, :delta, CURRENT_TIMESTAMP);
  1.212 END;
```

If multiple script files are specified, the averages are reported separately for each script file.

Note that collecting the additional timing information needed for per-statement latency computation adds some overhead. This will slow average execution speed and lower the computed TPS. The amount of slowdown varies significantly depending on platform and hardware. Comparing average TPS values with and without latency reporting enabled is a good way to measure if the timing overhead is significant.

**Good Practices**

It is very easy to use pgbench to produce completely meaningless numbers. Here are some guidelines to help you get useful results.

In the first place, *never* believe any test that runs for only a few seconds. Use the `-t` or `-T` option to make the run last at least a few minutes, so as to average out noise. In some cases you could need hours to get numbers that are reproducible. It’s a good idea to try the test run a few times, to find out if your numbers are reproducible or not.

For the default TPC-B-like test scenario, the initialization scale factor (`-s`) should be at least as large as the largest number of clients you intend to test (`-c`); else you’ll mostly be measuring update contention. There are only `-s` rows in the `pgbench_branches` table, and every transaction wants
to update one of them, so \(-c\) values in excess of \(-s\) will undoubtedly result in lots of transactions blocked waiting for other transactions.

The default test scenario is also quite sensitive to how long it’s been since the tables were initialized: accumulation of dead rows and dead space in the tables changes the results. To understand the results you must keep track of the total number of updates and when vacuuming happens. If autovacuum is enabled it can result in unpredictable changes in measured performance.

A limitation of pgbench is that it can itself become the bottleneck when trying to test a large number of client sessions. This can be alleviated by running pgbench on a different machine from the database server, although low network latency will be essential. It might even be useful to run several pgbench instances concurrently, on several client machines, against the same database server.

Security

If untrusted users have access to a database that has not adopted a secure schema usage pattern, do not run pgbench in that database. pgbench uses unqualified names and does not manipulate the search path.
pg_config

Name

pg_config — retrieve information about the installed version of PostgreSQL

Synopsis

pg_config [option...]

Description

The pg_config utility prints configuration parameters of the currently installed version of PostgreSQL. It is intended, for example, to be used by software packages that want to interface to PostgreSQL to facilitate finding the required header files and libraries.

Options

To use pg_config, supply one or more of the following options:

--bindir

Print the location of user executables. Use this, for example, to find the psql program. This is normally also the location where the pg_config program resides.

--docdir

Print the location of documentation files.

--htmldir

Print the location of HTML documentation files.

--includedir

Print the location of C header files of the client interfaces.

--pkgincludedir

Print the location of other C header files.

--includedir-server

Print the location of C header files for server programming.

--libdir

Print the location of object code libraries.

--pkglibdir

Print the location of dynamically loadable modules, or where the server would search for them. (Other architecture-dependent data files might also be installed in this directory.)

--localedir

Print the location of locale support files. (This will be an empty string if locale support was not configured when PostgreSQL was built.)
--mandir
    Print the location of manual pages.

--sharedir
    Print the location of architecture-independent support files.

--sysconfdir
    Print the location of system-wide configuration files.

--pgxs
    Print the location of extension makefiles.

--configure
    Print the options that were given to the configure script when PostgreSQL was configured for building. This can be used to reproduce the identical configuration, or to find out with what options a binary package was built. (Note however that binary packages often contain vendor-specific custom patches.) See also the examples below.

--cc
    Print the value of the CC variable that was used for building PostgreSQL. This shows the C compiler used.

--cppflags
    Print the value of the CPPFLAGS variable that was used for building PostgreSQL. This shows C compiler switches needed at preprocessing time (typically, -I switches).

--cflags
    Print the value of the CFLAGS variable that was used for building PostgreSQL. This shows C compiler switches.

--cflags_sl
    Print the value of the CFLAGS_SL variable that was used for building PostgreSQL. This shows extra C compiler switches used for building shared libraries.

--ldflags
    Print the value of the LDFLAGS variable that was used for building PostgreSQL. This shows linker switches.

--ldflags_ex
    Print the value of the LDFLAGS_EX variable that was used for building PostgreSQL. This shows linker switches used for building executables only.

--ldflags_sl
    Print the value of the LDFLAGS_SL variable that was used for building PostgreSQL. This shows linker switches used for building shared libraries only.

--libs
    Print the value of the LIBS variable that was used for building PostgreSQL. This normally contains -l switches for external libraries linked into PostgreSQL.

--version
    Print the version of PostgreSQL.
--help

Show help about pg_config command line arguments, and exit.

If more than one option is given, the information is printed in that order, one item per line. If no options are given, all available information is printed, with labels.

Notes

The options --docdir, --pkgincludedir, --localedir, --mandir, --sharedir, --sysconfdir, --cc, --cppflags, --cflags, --cflags_sl, --ldflags, --ldflags_sl, and --libs were added in PostgreSQL 8.1. The option --htmldir was added in PostgreSQL 8.4. The option --ldflags_ex was added in PostgreSQL 9.0.

Example

To reproduce the build configuration of the current PostgreSQL installation, run the following command:

eval ./configure 'pg_config --configure'

The output of pg_config --configure contains shell quotation marks so arguments with spaces are represented correctly. Therefore, using eval is required for proper results.
pg_dump

Name
pg_dump — extract a PostgreSQL database into a script file or other archive file

Synopsis
pg_dump [connection-option...][option...][dbname]

Description
pg_dump is a utility for backing up a PostgreSQL database. It makes consistent backups even if the database is being used concurrently. pg_dump does not block other users accessing the database (readers or writers).

pg_dump only dumps a single database. To backup global objects that are common to all databases in a cluster, such as roles and table spaces, use pg_dumpall.

Dumps can be output in script or archive file formats. Script dumps are plain-text files containing the SQL commands required to reconstruct the database to the state it was in at the time it was saved. To restore from such a script, feed it to psql. Script files can be used to reconstruct the database even on other machines and other architectures; with some modifications, even on other SQL database products.

The alternative archive file formats must be used with pg_restore to rebuild the database. They allow pg_restore to be selective about what is restored, or even to reorder the items prior to being restored. The archive file formats are designed to be portable across architectures.

When used with one of the archive file formats and combined with pg_restore, pg_dump provides a flexible archival and transfer mechanism. pg_dump can be used to backup an entire database, then pg_restore can be used to examine the archive and/or select which parts of the database are to be restored. The most flexible output file formats are the “custom” format (-Fc) and the “directory” format (-Fd). They allow for selection and reordering of all archived items, support parallel restoration, and are compressed by default. The “directory” format is the only format that supports parallel dumps.

While running pg_dump, one should examine the output for any warnings (printed on standard error), especially in light of the limitations listed below.

Options
The following command-line options control the content and format of the output.

dbname
  Specifies the name of the database to be dumped. If this is not specified, the environment variable PGDATABASE is used. If that is not set, the user name specified for the connection is used.
-a
--data-only

Dump only the data, not the schema (data definitions). Table data, large objects, and sequence values are dumped.

This option is similar to, but for historical reasons not identical to, specifying --section=data.

-b
--blobs

Include large objects in the dump. This is the default behavior except when --schema, --table, or --schema-only is specified. The -b switch is therefore only useful to add large objects to dumps where a specific schema or table has been requested. Note that blobs are considered data and therefore will be included when --data-only is used, but not when --schema-only is.

-c
--clean

Output commands to clean (drop) database objects prior to outputting the commands for creating them. (Unless --if-exists is also specified, restore might generate some harmless error messages, if any objects were not present in the destination database.)

This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call pg_restore.

-C
--create

Begin the output with a command to create the database itself and reconnect to the created database. (With a script of this form, it doesn’t matter which database in the destination installation you connect to before running the script.) If --clean is also specified, the script drops and recreates the target database before reconnecting to it.

This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call pg_restore.

-E encoding
--encoding=encoding

Create the dump in the specified character set encoding. By default, the dump is created in the database encoding. (Another way to get the same result is to set the PGCLIENTENCODING environment variable to the desired dump encoding.)

-f file
--file=file

Send output to the specified file. This parameter can be omitted for file based output formats, in which case the standard output is used. It must be given for the directory output format however, where it specifies the target directory instead of a file. In this case the directory is created by pg_dump and must not exist before.

-F format
--format=format

Selects the format of the output. format can be one of the following:

p
plain

Output a plain-text SQL script file (the default).
Output a custom-format archive suitable for input into pg_restore. Together with the directory output format, this is the most flexible output format in that it allows manual selection and reordering of archived items during restore. This format is also compressed by default.

Output a directory-format archive suitable for input into pg_restore. This will create a directory with one file for each table and blob being dumped, plus a so-called Table of Contents file describing the dumped objects in a machine-readable format that pg_restore can read. A directory format archive can be manipulated with standard Unix tools; for example, files in an uncompressed archive can be compressed with the gzip tool. This format is compressed by default and also supports parallel dumps.

Output a tar-format archive suitable for input into pg_restore. The tar format is compatible with the directory format: extracting a tar-format archive produces a valid directory-format archive. However, the tar format does not support compression. Also, when using tar format the relative order of table data items cannot be changed during restore.

Run the dump in parallel by dumping \texttt{njobs} tables simultaneously. This option reduces the time of the dump but it also increases the load on the database server. You can only use this option with the directory output format because this is the only output format where multiple processes can write their data at the same time.

pg_dump will open \texttt{njobs} + 1 connections to the database, so make sure your max_connections setting is high enough to accommodate all connections.

Requesting exclusive locks on database objects while running a parallel dump could cause the dump to fail. The reason is that the pg_dump master process requests shared locks on the objects that the worker processes are going to dump later in order to make sure that nobody deletes them and makes them go away while the dump is running. If another client then requests an exclusive lock on a table, that lock will not be granted but will be queued waiting for the shared lock of the master process to be released. Consequently any other access to the table will not be granted either and will queue after the exclusive lock request. This includes the worker process trying to dump the table. Without any precautions this would be a classic deadlock situation. To detect this conflict, the pg_dump worker process requests another shared lock using the NOWAIT option. If the worker process is not granted this shared lock, somebody else must have requested an exclusive lock in the meantime and there is no way to continue with the dump, so pg_dump has no choice but to abort the dump.

For a consistent backup, the database server needs to support synchronized snapshots, a feature that was introduced in PostgreSQL 9.2. With this feature, database clients can ensure they see the same data set even though they use different connections. pg_dump \texttt{-j} uses multiple database connections; it connects to the database once with the master process and once again for each worker job. Without the synchronized snapshot feature, the different worker jobs wouldn’t be guaranteed to see the same data in each connection, which could lead to an inconsistent backup.
If you want to run a parallel dump of a pre-9.2 server, you need to make sure that the database content doesn’t change from between the time the master connects to the database until the last worker job has connected to the database. The easiest way to do this is to halt any data modifying processes (DDL and DML) accessing the database before starting the backup. You also need to specify the --no-synchronized-snapshots parameter when running `pg_dump -j` against a pre-9.2 PostgreSQL server.

```bash
-n schema
--schema=schema
```

Dump only schemas matching `schema`; this selects both the schema itself, and all its contained objects. When this option is not specified, all non-system schemas in the target database will be dumped. Multiple schemas can be selected by writing multiple `-n` switches. Also, the `schema` parameter is interpreted as a pattern according to the same rules used by `psql`'s `\d` commands (see Patterns), so multiple schemas can also be selected by writing wildcard characters in the pattern. When using wildcards, be careful to quote the pattern if needed to prevent the shell from expanding the wildcards; see Examples.

**Note:** When `-n` is specified, `pg_dump` makes no attempt to dump any other database objects that the selected schema(s) might depend upon. Therefore, there is no guarantee that the results of a specific-schema dump can be successfully restored by themselves into a clean database.

**Note:** Non-schema objects such as blobs are not dumped when `-n` is specified. You can add blobs back to the dump with the `--blobs` switch.

```bash
-N schema
--exclude-schema=schema
```

Do not dump any schemas matching the `schema` pattern. The pattern is interpreted according to the same rules as for `-n`. `-N` can be given more than once to exclude schemas matching any of several patterns.

When both `-n` and `-N` are given, the behavior is to dump just the schemas that match at least one `-n` switch but no `-N` switches. If `-N` appears without `-n`, then schemas matching `-N` are excluded from what is otherwise a normal dump.

```bash
-o
--oids
```

Dump object identifiers (OIDs) as part of the data for every table. Use this option if your application references the OID columns in some way (e.g., in a foreign key constraint). Otherwise, this option should not be used.

```bash
-O
--no-owner
```

Do not output commands to set ownership of objects to match the original database. By default, `pg_dump` issues `ALTER OWNER` or `SET SESSION AUTHORIZATION` statements to set ownership of created database objects. These statements will fail when the script is run unless it is started by a superuser (or the same user that owns all of the objects in the script). To make a script that can be restored by any user, but will give that user ownership of all the objects, specify `-O`.

This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call `pg_restore`. 

1819
-R
--no-reconnect

This option is obsolete but still accepted for backwards compatibility.

-s
--schema-only

Dump only the object definitions (schema), not data.

This option is the inverse of --data-only. It is similar to, but for historical reasons not identical to, specifying --section=pre-data --section=post-data.

(Do not confuse this with the --schema option, which uses the word “schema” in a different meaning.)

To exclude table data for only a subset of tables in the database, see --exclude-table-data.

-S username
--superuser=username

Specify the superuser user name to use when disabling triggers. This is relevant only if --disable-triggers is used. (Usually, it’s better to leave this out, and instead start the resulting script as superuser.)

-t table
--table=table

Dump only tables with names matching table. For this purpose, “table” includes views, materialized views, sequences, and foreign tables. Multiple tables can be selected by writing multiple -t switches. Also, the table parameter is interpreted as a pattern according to the same rules used by psql’s \d commands (see Patterns), so multiple tables can also be selected by writing wildcard characters in the pattern. When using wildcards, be careful to quote the pattern if needed to prevent the shell from expanding the wildcards; see Examples.

The -n and -N switches have no effect when -t is used, because tables selected by -t will be dumped regardless of those switches, and non-table objects will not be dumped.

Note: When -t is specified, pg_dump makes no attempt to dump any other database objects that the selected table(s) might depend upon. Therefore, there is no guarantee that the results of a specific-table dump can be successfully restored by themselves into a clean database.

Note: The behavior of the -t switch is not entirely upward compatible with pre-8.2 PostgreSQL versions. Formerly, writing -t tab would dump all tables named tab, but now it just dumps whichever one is visible in your default search path. To get the old behavior you can write -t ‘*.tab’. Also, you must write something like -t sch.tab to select a table in a particular schema, rather than the old locution of -n sch -t tab.

-T table
--exclude-table=table

Do not dump any tables matching the table pattern. The pattern is interpreted according to the same rules as for -t. -T can be given more than once to exclude tables matching any of several patterns.
When both -t and -T are given, the behavior is to dump just the tables that match at least one -t switch but no -T switches. If -T appears without -t, then tables matching -T are excluded from what is otherwise a normal dump.

-v
--verbose

Specifies verbose mode. This will cause pg_dump to output detailed object comments and start/stop times to the dump file, and progress messages to standard error.

-V
--version

Print the pg_dump version and exit.

-x
--no-privileges
--no-acl

Prevent dumping of access privileges (grant/revoke commands).

-Z 0..9
--compress=0..9

Specify the compression level to use. Zero means no compression. For the custom archive format, this specifies compression of individual table-data segments, and the default is to compress at a moderate level. For plain text output, setting a nonzero compression level causes the entire output file to be compressed, as though it had been fed through gzip; but the default is not to compress. The tar archive format currently does not support compression at all.

--binary-upgrade

This option is for use by in-place upgrade utilities. Its use for other purposes is not recommended or supported. The behavior of the option may change in future releases without notice.

--column-inserts
--attribute-inserts

Dump data as INSERT commands with explicit column names (INSERT INTO table (column, ...) VALUES ...). This will make restoration very slow; it is mainly useful for making dumps that can be loaded into non-PostgreSQL databases. However, since this option generates a separate command for each row, an error in reloading a row causes only that row to be lost rather than the entire table contents.

--disable-dollar-quoting

This option disables the use of dollar quoting for function bodies, and forces them to be quoted using SQL standard string syntax.

--disable-triggers

This option is relevant only when creating a data-only dump. It instructs pg_dump to include commands to temporarily disable triggers on the target tables while the data is reloaded. Use this if you have referential integrity checks or other triggers on the tables that you do not want to invoke during data reload. Presently, the commands emitted for --disable-triggers must be done as superuser. So, you should also specify a superuser name with -S, or preferably be careful to start the resulting script as a superuser.

This option is only meaningful for the plain-text format. For the archive formats, you can specify the option when you call pg_restore.
---enable-row-security
This option is relevant only when dumping the contents of a table which has row security. By
default, pg_dump will set row_security to off, to ensure that all data is dumped from the table. If
the user does not have sufficient privileges to bypass row security, then an error is thrown. This
parameter instructs pg_dump to set row_security to on instead, allowing the user to dump the
parts of the contents of the table that they have access to.
Note that if you use this option currently, you probably also want the dump be in INSERT format,
as the COPY FROM during restore does not support row security.

---exclude-table-data=table
Do not dump data for any tables matching the table pattern. The pattern is interpreted according
to the same rules as for -t. --exclude-table-data can be given more than once to exclude
tables matching any of several patterns. This option is useful when you need the definition of a
particular table even though you do not need the data in it.
To exclude data for all tables in the database, see --schema-only.

---if-exists
Use conditional commands (i.e. add an IF EXISTS clause) when cleaning database objects. This
option is not valid unless --clean is also specified.

---inserts
Dump data as INSERT commands (rather than COPY). This will make restoration very slow; it is
mainly useful for making dumps that can be loaded into non-PostgreSQL databases. However,
since this option generates a separate command for each row, an error in reloading a row causes
only that row to be lost rather than the entire table contents. Note that the restore might fail
altogether if you have rearranged column order. The --column-inserts option is safe against
column order changes, though even slower.

---lock-wait-timeout=timeout
Do not wait forever to acquire shared table locks at the beginning of the dump. Instead fail if
unable to lock a table within the specified timeout. The timeout may be specified in any of the
formats accepted by SET statement_timeout. (Allowed values vary depending on the server
version you are dumping from, but an integer number of milliseconds is accepted by all versions
since 7.3. This option is ignored when dumping from a pre-7.3 server.)

---no-security-labels
Do not dump security labels.

---no-synchronized-snapshots
This option allows running pg_dump -j against a pre-9.2 server, see the documentation of the
-j parameter for more details.

---no-tablespaces
Do not output commands to select tablespaces. With this option, all objects will be created in
whichever tablespace is the default during restore.
This option is only meaningful for the plain-text format. For the archive formats, you can specify
the option when you call pg_restore.

---no-unlogged-table-data
Do not dump the contents of unlogged tables. This option has no effect on whether or not the table
definitions (schema) are dumped; it only suppresses dumping the table data. Data in unlogged
tables is always excluded when dumping from a standby server.
--quote-all-identifiers

Force quoting of all identifiers. This option is recommended when dumping a database from a server whose PostgreSQL major version is different from pg_dump’s, or when the output is intended to be loaded into a server of a different major version. By default, pg_dump quotes only identifiers that are reserved words in its own major version. This sometimes results in compatibility issues when dealing with servers of other versions that may have slightly different sets of reserved words. Using --quote-all-identifiers prevents such issues, at the price of a harder-to-read dump script.

--section=sectionname

Only dump the named section. The section name can be pre-data, data, or post-data. This option can be specified more than once to select multiple sections. The default is to dump all sections.

The data section contains actual table data, large-object contents, and sequence values. Post-data items include definitions of indexes, triggers, rules, and constraints other than validated check constraints. Pre-data items include all other data definition items.

--serializable-deferrable

Use a serializable transaction for the dump, to ensure that the snapshot used is consistent with later database states; but do this by waiting for a point in the transaction stream at which no anomalies can be present, so that there isn’t a risk of the dump failing or causing other transactions to roll back with a serialization_failure. See Chapter 13 for more information about transaction isolation and concurrency control.

This option is not beneficial for a dump which is intended only for disaster recovery. It could be useful for a dump used to load a copy of the database for reporting or other read-only load sharing while the original database continues to be updated. Without it the dump may reflect a state which is not consistent with any serial execution of the transactions eventually committed. For example, if batch processing techniques are used, a batch may show as closed in the dump without all of the items which are in the batch appearing.

This option will make no difference if there are no read-write transactions active when pg_dump is started. If read-write transactions are active, the start of the dump may be delayed for an indeterminate length of time. Once running, performance with or without the switch is the same.

--snapshot=snapshotname

Use the specified synchronized snapshot when making a dump of the database (see Table 9-81 for more details).

This option is useful when needing to synchronize the dump with a logical replication slot (see Chapter 47) or with a concurrent session.

In the case of a parallel dump, the snapshot name defined by this option is used rather than taking a new snapshot.

--strict-names

Require that each schema (-n/--schema) and table (-t/--table) qualifier match at least one schema/table in the database to be dumped. Note that if none of the schema/table qualifiers find matches, pg_dump will generate an error even without --strict-names.

This option has no effect on -N/--exclude-schema, -T/--exclude-table, or --exclude-table-data. An exclude pattern failing to match any objects is not considered an error.
pg_dump

--use-set-session-authorization

Output SQL-standard SET SESSION AUTHORIZATION commands instead of ALTER OWNER commands to determine object ownership. This makes the dump more standards-compatible, but depending on the history of the objects in the dump, might not restore properly. Also, a dump using SET SESSION AUTHORIZATION will certainly require superuser privileges to restore correctly, whereas ALTER OWNER requires lesser privileges.

-?
--help

Show help about pg_dump command line arguments, and exit.

The following command-line options control the database connection parameters.

-d dbname
--dbname=
dbname

Specifies the name of the database to connect to. This is equivalent to specifying dbname as the first non-option argument on the command line.

If this parameter contains an = sign or starts with a valid URI prefix (postgresql:// or postgres://), it is treated as a conninfo string. See Section 32.1 for more information.

-h host
--host=
host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

-p port
--port=
port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

-U username
--username=
username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force pg_dump to prompt for a password before connecting to a database.

This option is never essential, since pg_dump will automatically prompt for a password if the server demands password authentication. However, pg_dump will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.
--role=rolename

Specifies a role name to be used to create the dump. This option causes pg_dump to issue a SET ROLE rolename command after connecting to the database. It is useful when the authenticated user (specified by -U) lacks privileges needed by pg_dump, but can switch to a role with the required rights. Some installations have a policy against logging in directly as a superuser, and use of this option allows dumps to be made without violating the policy.

Environment

PGDATABASE
PGHOST
PGOPTIONS
PGPORT
PGUSER

Default connection parameters.

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).

Diagnostics

pg_dump internally executes SELECT statements. If you have problems running pg_dump, make sure you are able to select information from the database using, for example, psql. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

The database activity of pg_dump is normally collected by the statistics collector. If this is undesirable, you can set parameter track_counts to false via PGOPTIONS or the ALTER USER command.

Notes

If your database cluster has any local additions to the template1 database, be careful to restore the output of pg_dump into a truly empty database; otherwise you are likely to get errors due to duplicate definitions of the added objects. To make an empty database without any local additions, copy from template0 not template1, for example:

CREATE DATABASE foo WITH TEMPLATE template0;

When a data-only dump is chosen and the option --disable-triggers is used, pg_dump emits commands to disable triggers on user tables before inserting the data, and then commands to re-enable them after the data has been inserted. If the restore is stopped in the middle, the system catalogs might be left in the wrong state.

The dump file produced by pg_dump does not contain the statistics used by the optimizer to make query planning decisions. Therefore, it is wise to run ANALYZE after restoring from a dump file to ensure optimal performance; see Section 24.1.3 and Section 24.1.6 for more information. The dump file also does not contain any ALTER DATABASE ... SET commands; these settings are dumped by pg_dumpall, along with database users and other installation-wide settings.
Because pg_dump is used to transfer data to newer versions of PostgreSQL, the output of pg_dump can be expected to load into PostgreSQL server versions newer than pg_dump’s version. pg_dump can also dump from PostgreSQL servers older than its own version. (Currently, servers back to version 7.0 are supported.) However, pg_dump cannot dump from PostgreSQL servers newer than its own major version; it will refuse to even try, rather than risk making an invalid dump. Also, it is not guaranteed that pg_dump’s output can be loaded into a server of an older major version — not even if the dump was taken from a server of that version. Loading a dump file into an older server may require manual editing of the dump file to remove syntax not understood by the older server. Use of the --quote-all-identifiers option is recommended in cross-version cases, as it can prevent problems arising from varying reserved-word lists in different PostgreSQL versions.

**Examples**

To dump a database called mydb into a SQL-script file:

```
$ pg_dump mydb > db.sql
```

To reload such a script into a (freshly created) database named newdb:

```
$ psql -d newdb -f db.sql
```

To dump a database into a custom-format archive file:

```
$ pg_dump -Fc mydb > db.dump
```

To dump a database into a directory-format archive:

```
$ pg_dump -Fd mydb -f dumpdir
```

To dump a database into a directory-format archive in parallel with 5 worker jobs:

```
$ pg_dump -Fd mydb -j 5 -f dumpdir
```

To reload an archive file into a (freshly created) database named newdb:

```
$ pg_restore -d newdb db.dump
```

To dump a single table named mytab:

```
$ pg_dump -t mytab mydb > db.sql
```

To dump all tables whose names start with emp in the detroit schema, except for the table named employee_log:

```
$ pg_dump -t 'detroit.emp*' -T detroit.employee_log mydb > db.sql
```
To dump all schemas whose names start with east or west and end in gsm, excluding any schemas whose names contain the word test:

```
$ pg_dump -n 'east*gsm' -n 'west*gsm' -N '*test*' mydb > db.sql
```

The same, using regular expression notation to consolidate the switches:

```
$ pg_dump -n '(east|west)*gsm -N '*test*' mydb > db.sql
```

To dump all database objects except for tables whose names begin with ts_:

```
$ pg_dump -T 'ts_*' mydb > db.sql
```

To specify an upper-case or mixed-case name in -t and related switches, you need to double-quote the name; else it will be folded to lower case (see Patterns). But double quotes are special to the shell, so in turn they must be quoted. Thus, to dump a single table with a mixed-case name, you need something like

```
$ pg_dump -t "\"MixedCaseName\"" mydb > mytab.sql
```

See Also

pg_dumpall, pg_restore, psql
**pg_dumpall**

**Name**

`pg_dumpall` — extract a PostgreSQL database cluster into a script file

**Synopsis**

```
pg_dumpall [connection-option...][option...]
```

**Description**

`pg_dumpall` is a utility for writing out (“dumping”) all PostgreSQL databases of a cluster into one script file. The script file contains SQL commands that can be used as input to `psql` to restore the databases. It does this by calling `pg_dump` for each database in a cluster. `pg_dumpall` also dumps global objects that are common to all databases. (`pg_dump` does not save these objects.) This currently includes information about database users and groups, tablespaces, and properties such as access permissions that apply to databases as a whole.

Since `pg_dumpall` reads tables from all databases you will most likely have to connect as a database superuser in order to produce a complete dump. Also you will need superuser privileges to execute the saved script in order to be allowed to add users and groups, and to create databases.

The SQL script will be written to the standard output. Use the `-f/--file` option or shell operators to redirect it into a file.

`pg_dumpall` needs to connect several times to the PostgreSQL server (once per database). If you use password authentication it will ask for a password each time. It is convenient to have a `~/.pgpass` file in such cases. See Section 32.15 for more information.

**Options**

The following command-line options control the content and format of the output.

- `-a`
  `--data-only`
  Dump only the data, not the schema (data definitions).

- `-c`
  `--clean`
  Include SQL commands to clean (drop) databases before recreating them. DROP commands for roles and tablespaces are added as well.

- `-f filename`
  `--file=filename`
  Send output to the specified file. If this is omitted, the standard output is used.
pg_dumpall

-g
--globals-only

  Dump only global objects (roles and tablespaces), no databases.

-o
--oids

  Dump object identifiers (OIDs) as part of the data for every table. Use this option if your application references the OID columns in some way (e.g., in a foreign key constraint). Otherwise, this option should not be used.

-O
--no-owner

  Do not output commands to set ownership of objects to match the original database. By default, pg_dumpall issues ALTER OWNER or SET SESSION AUTHORIZATION statements to set ownership of created schema elements. These statements will fail when the script is run unless it is started by a superuser (or the same user that owns all of the objects in the script). To make a script that can be restored by any user, but will give that user ownership of all the objects, specify -O.

-r
--roles-only

  Dump only roles, no databases or tablespaces.

-s
--schema-only

  Dump only the object definitions (schema), not data.

-S username
--superuser=username

  Specify the superuser user name to use when disabling triggers. This is relevant only if --disable-triggers is used. (Usually, it’s better to leave this out, and instead start the resulting script as superuser.)

-t
--tablespaces-only

  Dump only tablespaces, no databases or roles.

-v
--verbose

  Specifies verbose mode. This will cause pg_dumpall to output start/stop times to the dump file, and progress messages to standard error. It will also enable verbose output in pg_dump.

-V
--version

  Print the pg_dumpall version and exit.

-x
--no-privileges
--no-acl

  Prevent dumping of access privileges (grant/revoke commands).
--binary-upgrade

This option is for use by in-place upgrade utilities. Its use for other purposes is not recommended or supported. The behavior of the option may change in future releases without notice.

--column-inserts  --attribute-inserts

Dump data as INSERT commands with explicit column names (INSERT INTO table (column, ...) VALUES ...). This will make restoration very slow; it is mainly useful for making dumps that can be loaded into non-PostgreSQL databases.

--disable-dollar-quoting

This option disables the use of dollar quoting for function bodies, and forces them to be quoted using SQL standard string syntax.

--disable-triggers

This option is relevant only when creating a data-only dump. It instructs pg_dumpall to include commands to temporarily disable triggers on the target tables while the data is reloaded. Use this if you have referential integrity checks or other triggers on the tables that you do not want to invoke during data reload.

Presently, the commands emitted for --disable-triggers must be done as superuser. So, you should also specify a superuser name with -S, or preferably be careful to start the resulting script as a superuser.

--if-exists

Use conditional commands (i.e. add an IF EXISTS clause) to clean databases and other objects. This option is not valid unless --clean is also specified.

--inserts

Dump data as INSERT commands (rather than COPY). This will make restoration very slow; it is mainly useful for making dumps that can be loaded into non-PostgreSQL databases. Note that the restore might fail altogether if you have rearranged column order. The --column-inserts option is safer, though even slower.

--lock-wait-timeout=timeout

Do not wait forever to acquire shared table locks at the beginning of the dump. Instead, fail if unable to lock a table within the specified timeout. The timeout may be specified in any of the formats accepted by SET statement_timeout. Allowed values vary depending on the server version you are dumping from, but an integer number of milliseconds is accepted by all versions since 7.3. This option is ignored when dumping from a pre-7.3 server.

--no-security-labels

Do not dump security labels.

--no-tablespaces

Do not output commands to create tablespaces nor select tablespaces for objects. With this option, all objects will be created in whichever tablespace is the default during restore.

--no-unlogged-table-data

Do not dump the contents of unlogged tables. This option has no effect on whether or not the table definitions (schema) are dumped; it only suppresses dumping the table data.
pg_dumpall

--quote-all-identifiers

Force quoting of all identifiers. This option is recommended when dumping a database from a server whose PostgreSQL major version is different from pg_dumpall’s, or when the output is intended to be loaded into a server of a different major version. By default, pg_dumpall quotes only identifiers that are reserved words in its own major version. This sometimes results in compatibility issues when dealing with servers of other versions that may have slightly different sets of reserved words. Using --quote-all-identifiers prevents such issues, at the price of a harder-to-read dump script.

--use-set-session-authorization

Output SQL-standard SET SESSION AUTHORIZATION commands instead of ALTER OWNER commands to determine object ownership. This makes the dump more standards compatible, but depending on the history of the objects in the dump, might not restore properly.

-?
--help

Show help about pg_dumpall command line arguments, and exit.

The following command-line options control the database connection parameters.

-d connstr
--dbname=connstr

Specifies parameters used to connect to the server, as a connection string. See Section 32.1.1 for more information.

The option is called --dbname for consistency with other client applications, but because pg_dumpall needs to connect to many databases, database name in the connection string will be ignored. Use -l option to specify the name of the database used to dump global objects and to discover what other databases should be dumped.

-h host
--host=host

Specifies the host name of the machine on which the database server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

-l dbname
--database=dbname

Specifies the name of the database to connect to for dumping global objects and discovering what other databases should be dumped. If not specified, the postgres database will be used, and if that does not exist, template1 will be used.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

-U username
--username=username

User name to connect as.
---

**pg_dumpall**

```
-\w
--no-password

Never issue a password prompt. If the server requires password authentication and a password
is not available by other means such as a .pgpass file, the connection attempt will fail. This
option can be useful in batch jobs and scripts where no user is present to enter a password.
```

```
-\W
--password

Force pg_dumpall to prompt for a password before connecting to a database.
This option is never essential, since pg_dumpall will automatically prompt for a password if the
server demands password authentication. However, pg_dumpall will waste a connection attempt
finding out that the server wants a password. In some cases it is worth typing -w to avoid the
extra connection attempt.

Note that the password prompt will occur again for each database to be dumped. Usually, it’s
better to set up a ~/.pgpass file than to rely on manual password entry.
```

```
--role=rolename

Specifies a role name to be used to create the dump. This option causes pg_dumpall to issue a SET
ROLE rolename command after connecting to the database. It is useful when the authenticated
user (specified by -u) lacks privileges needed by pg_dumpall, but can switch to a role with the
required rights. Some installations have a policy against logging in directly as a superuser, and
use of this option allows dumps to be made without violating the policy.
```

### Environment

**PGHOST**
**PGOPTIONS**
**PGPORT**
**PGUSER**

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by
libpq (see Section 32.14).

### Notes

Since pg_dumpall calls pg_dump internally, some diagnostic messages will refer to pg_dump.

Once restored, it is wise to run **ANALYZE** on each database so the optimizer has useful statistics. You
can also run vacuumdb -a -z to analyze all databases.

pg_dumpall requires all needed tablespace directories to exist before the restore; otherwise, database
creation will fail for databases in non-default locations.
Examples

To dump all databases:

$ pg_dumpall > db.out

To reload database(s) from this file, you can use:

$ psql -f db.out postgres

(It is not important to which database you connect here since the script file created by pg_dumpall will contain the appropriate commands to create and connect to the saved databases.)

See Also

Check pg_dump for details on possible error conditions.
**pg_isready**

**Name**

pg_isready — check the connection status of a PostgreSQL server

**Synopsis**

pg_isready [connection-option...][option...]

**Description**

pg_isready is a utility for checking the connection status of a PostgreSQL database server. The exit status specifies the result of the connection check.

**Options**

- `-d dbname`
  `-dbname=`
  Specifies the name of the database to connect to.
  If this parameter contains an `=` sign or starts with a valid URI prefix (`postgresql://` or `postgres://`), it is treated as a conninfo string. See Section 32.1.1 for more information.

- `-h hostname`
  `-host=`
  Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix-domain socket.

- `-p port`
  `-port=`
  Specifies the TCP port or the local Unix-domain socket file extension on which the server is listening for connections. Defaults to the value of the `PGPORT` environment variable or, if not set, to the port specified at compile time, usually 5432.

- `-q`
  `-quiet`
  Do not display status message. This is useful when scripting.

- `-t seconds`
  `-timeout=`
  The maximum number of seconds to wait when attempting connection before returning that the server is not responding. Setting to 0 disables. The default is 3 seconds.

- `-U username`
  `-username=`
  Connect to the database as the user `username` instead of the default.
-V
--version

Print the pg_isready version and exit.

-?
--help

Show help about pg_isready command line arguments, and exit.

Exit Status

pg_isready returns 0 to the shell if the server is accepting connections normally, 1 if the server is rejecting connections (for example during startup), 2 if there was no response to the connection attempt, and 3 if no attempt was made (for example due to invalid parameters).

Environment

pg_isready, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).

Notes

It is not necessary to supply correct user name, password, or database name values to obtain the server status; however, if incorrect values are provided, the server will log a failed connection attempt.

Examples

Standard Usage:

$ pg_isready /tmp:5432 - accepting connections
$ echo $? 0

Running with connection parameters to a PostgreSQL cluster in startup:

$ pg_isready -h localhost -p 5433
localhost:5433 - rejecting connections
$ echo $? 1

Running with connection parameters to a non-responsive PostgreSQL cluster:

$ pg_isready -h someremotehost
someremotehost:5432 - no response
$ echo $? 2
pg_isready
\textbf{pg\_receivexlog}

\section*{Name}

pg\_receivexlog — stream transaction logs from a PostgreSQL server

\section*{Synopsis}

pg\_receivexlog[option...]

\section*{Description}

pg\_receivexlog is used to stream the transaction log from a running PostgreSQL cluster. The transaction log is streamed using the streaming replication protocol, and is written to a local directory of files. This directory can be used as the archive location for doing a restore using point-in-time recovery (see Section 25.3).

pg\_receivexlog streams the transaction log in real time as it’s being generated on the server, and does not wait for segments to complete like archive\_command does. For this reason, it is not necessary to set archive\_timeout when using pg\_receivexlog.

Unlike the WAL receiver of a PostgreSQL standby server, pg\_receivexlog by default flushes WAL data only when a WAL file is closed. The option \texttt{--synchronous} must be specified to flush WAL data in real time. Since pg\_receivexlog does not apply WAL, you should not allow it to become a synchronous standby when synchronous\_commit equals remote\_apply. If it does, it will appear to be a standby that never catches up, and will cause transaction commits to block. To avoid this, you should either configure an appropriate value for synchronous\_standby\_names, or specify application\_name for pg\_receivexlog that does not match it, or change the value of synchronous\_commit to something other than remote\_apply.

The transaction log is streamed over a regular PostgreSQL connection and uses the replication protocol. The connection must be made with a superuser or a user having \texttt{REPLICATION} permissions (see Section 21.2), and \texttt{pg\_hba.conf} must permit the replication connection. The server must also be configured with max\_wal\_senders set high enough to leave at least one session available for the stream.

If the connection is lost, or if it cannot be initially established, with a non-fatal error, pg\_receivexlog will retry the connection indefinitely, and reestablish streaming as soon as possible. To avoid this behavior, use the \texttt{--n} parameter.

\section*{Options}

\begin{description}
\item \texttt{-D directory}
\item \texttt{--directory=directory}
\end{description}

Directory to write the output to.

This parameter is required.
pg_receivexlog

--if-not-exists
Do not error out when --create-slot is specified and a slot with the specified name already exists.

-n
--no-loop
Don’t loop on connection errors. Instead, exit right away with an error.

-s interval
--status-interval=interval
Specifies the number of seconds between status packets sent back to the server. This allows for easier monitoring of the progress from server. A value of zero disables the periodic status updates completely, although an update will still be sent when requested by the server, to avoid timeout disconnect. The default value is 10 seconds.

-S slotname
--slot=slotname
Require pg_receivexlog to use an existing replication slot (see Section 26.2.6). When this option is used, pg_receivexlog will report a flush position to the server, indicating when each segment has been synchronized to disk so that the server can remove that segment if it is not otherwise needed.

When the replication client of pg_receivexlog is configured on the server as a synchronous standby, then using a replication slot will report the flush position to the server, but only when a WAL file is closed. Therefore, that configuration will cause transactions on the primary to wait for a long time and effectively not work satisfactorily. The option --synchronous (see below) must be specified in addition to make this work correctly.

--synchronous
Flush the WAL data to disk immediately after it has been received. Also send a status packet back to the server immediately after flushing, regardless of --status-interval.

This option should be specified if the replication client of pg_receivexlog is configured on the server as a synchronous standby, to ensure that timely feedback is sent to the server.

-v
--verbose
Enables verbose mode.

The following command-line options control the database connection parameters.

-d connstr
--dbname=connstr
Specifies parameters used to connect to the server, as a connection string. See Section 32.1.1 for more information.

The option is called --dbname for consistency with other client applications, but because pg_receivexlog doesn’t connect to any particular database in the cluster, database name in the connection string will be ignored.

-h host
--host=host
Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.
pg_receivexlog

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the $PGPORT$ environment variable, if set, or a compiled-in default.

-U username
--username=..

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force pg_receivexlog to prompt for a password before connecting to a database.

This option is never essential, since pg_receivexlog will automatically prompt for a password if the server demands password authentication. However, pg_receivexlog will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

pg_receivexlog can perform one of the two following actions in order to control physical replication slots:

--create-slot

Create a new physical replication slot with the name specified in --slot, then exit.

--drop-slot

Drop the replication slot with the name specified in --slot, then exit.

Other options are also available:

-V
--version

Print the pg_receivexlog version and exit.

-?
--help

Show help about pg_receivexlog command line arguments, and exit.

Environment

This utility, like most other PostgreSQL utilities, uses the environment variables supported by libpq (see Section 32.14).
Notes

When using `pg_receivexlog` instead of `archive_command` as the main WAL backup method, it is strongly recommended to use replication slots. Otherwise, the server is free to recycle or remove transaction log files before they are backed up, because it does not have any information, either from `archive_command` or the replication slots, about how far the WAL stream has been archived. Note, however, that a replication slot will fill up the server’s disk space if the receiver does not keep up with fetching the WAL data.

Examples

To stream the transaction log from the server at `mydbserver` and store it in the local directory `/usr/local/pgsql/archive`:

```
$ pg_receivexlog -h mydbserver -D /usr/local/pgsql/archive
```

See Also

`pg_basebackup`
pg_recvlogical

Name
pg_recvlogical — control PostgreSQL logical decoding streams

Synopsis
pg_recvlogical [option...]

Description
pg_recvlogical controls logical decoding replication slots and streams data from such replication slots.
It creates a replication-mode connection, so it is subject to the same constraints as pg_recvexlog, plus those for logical replication (see Chapter 47).

Options
At least one of the following options must be specified to select an action:

--create-slot
Create a new logical replication slot with the name specified by --slot, using the output plugin specified by --plugin, for the database specified by --dbname.

--drop-slot
Drop the replication slot with the name specified by --slot, then exit.

--start
Begin streaming changes from the logical replication slot specified by --slot, continuing until terminated by a signal. If the server side change stream ends with a server shutdown or discon-nect, retry in a loop unless --no-loop is specified.
The stream format is determined by the output plugin specified when the slot was created.
The connection must be to the same database used to create the slot.

--create-slot and --start can be specified together. --drop-slot cannot be combined with another action.

The following command-line options control the location and format of the output and other replication behavior:

-f filename
--file=filename
Write received and decoded transaction data into this file. Use - for stdout.
Specifies how often pg_recvlogical should issue \texttt{fsync()} calls to ensure the output file is safely flushed to disk.

The server will occasionally request the client to perform a flush and report the flush position to the server. This setting is in addition to that, to perform flushes more frequently.

Specifying an interval of 0 disables issuing \texttt{fsync()} calls altogether, while still reporting progress to the server. In this case, data could be lost in the event of a crash.

In \texttt{--start} mode, start replication from the given LSN. For details on the effect of this, see the documentation in Chapter 47 and Section 51.3. Ignored in other modes.

Do not error out when \texttt{--create-slot} is specified and a slot with the specified name already exists.

When the connection to the server is lost, do not retry in a loop, just exit.

Pass the option \texttt{name} to the output plugin with, if specified, the option value \texttt{value}. Which options exist and their effects depends on the used output plugin.

When creating a slot, use the specified logical decoding output plugin. See Chapter 47. This option has no effect if the slot already exists.

This option has the same effect as the option of the same name in pg_recvlogical. See the description there.

In \texttt{--start} mode, use the existing logical replication slot named \texttt{slot_name}. In \texttt{--create-slot} mode, create the slot with this name. In \texttt{--drop-slot} mode, delete the slot with this name.

Enables verbose mode.

The following command-line options control the database connection parameters.
The database to connect to. See the description of the actions for what this means in detail. This can be a libpq connection string; see Section 32.1.1 for more information. Defaults to user name.

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

User name to connect as. Defaults to current operating system user name.

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

Force pg_recvlogical to prompt for a password before connecting to a database.

This option is never essential, since pg_recvlogical will automatically prompt for a password if the server demands password authentication. However, pg_recvlogical will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing \texttt{-W} to avoid the extra connection attempt.

The following additional options are available:

Print the pg_recvlogical version and exit.

Show help about pg_recvlogical command line arguments, and exit.

Environment

This utility, like most other PostgreSQL utilities, uses the environment variables supported by libpq (see Section 32.14).
Examples

See Section 47.1 for an example.

See Also

pg_receivexlog
**pg_restore**

**Name**

`pg_restore` — restore a PostgreSQL database from an archive file created by `pg_dump`

**Synopsis**

```
pg_restore [connection-option...] [option...] [filename]
```

**Description**

`pg_restore` is a utility for restoring a PostgreSQL database from an archive created by `pg_dump` in one of the non-plain-text formats. It will issue the commands necessary to reconstruct the database to the state it was in at the time it was saved. The archive files also allow `pg_restore` to be selective about what is restored, or even to reorder the items prior to being restored. The archive files are designed to be portable across architectures.

`pg_restore` can operate in two modes. If a database name is specified, `pg_restore` connects to that database and restores archive contents directly into the database. Otherwise, a script containing the SQL commands necessary to rebuild the database is created and written to a file or standard output. This script output is equivalent to the plain text output format of `pg_dump`. Some of the options controlling the output are therefore analogous to `pg_dump` options.

Obviously, `pg_restore` cannot restore information that is not present in the archive file. For instance, if the archive was made using the “dump data as `INSERT` commands” option, `pg_restore` will not be able to load the data using `COPY` statements.

**Options**

`pg_restore` accepts the following command line arguments.

- `filename`
  Specifies the location of the archive file (or directory, for a directory-format archive) to be restored. If not specified, the standard input is used.

- `-a`  
  `--data-only`
  Restore only the data, not the schema (data definitions). Table data, large objects, and sequence values are restored, if present in the archive.

  This option is similar to, but for historical reasons not identical to, specifying `--section=data`.

- `-c`  
  `--clean`
  Clean (drop) database objects before recreating them. (Unless `--if-exists` is used, this might generate some harmless error messages, if any objects were not present in the destination database.)

1845
pg_restore

-C
--create

Create the database before restoring into it. If --clean is also specified, drop and recreate the target database before connecting to it.

When this option is used, the database named with -d is used only to issue the initial DROP DATABASE and CREATE DATABASE commands. All data is restored into the database name that appears in the archive.

-d dbname
--dbname=dbname

Connect to database dbname and restore directly into the database.

-e
--exit-on-error

Exit if an error is encountered while sending SQL commands to the database. The default is to continue and to display a count of errors at the end of the restoration.

-f filename
--file=filename

Specify output file for generated script, or for the listing when used with -l. Use - for the standard output, which is also the default.

-F format
--format=format

Specify format of the archive. It is not necessary to specify the format, since pg_restore will determine the format automatically. If specified, it can be one of the following:

c
  custom

  The archive is in the custom format of pg_dump.

d
directory

  The archive is a directory archive.

t
tar

  The archive is a tar archive.

-I index
--index=index

Restore definition of named index only. Multiple indexes may be specified with multiple -I switches.

-\( number-of-jobs\)
--jobs=number-of-jobs

Run the most time-consuming parts of pg_restore — those which load data, create indexes, or create constraints — using multiple concurrent jobs. This option can dramatically reduce the time to restore a large database to a server running on a multiprocessor machine.

Each job is one process or one thread, depending on the operating system, and uses a separate connection to the server.
The optimal value for this option depends on the hardware setup of the server, of the client, and of the network. Factors include the number of CPU cores and the disk setup. A good place to start is the number of CPU cores on the server, but values larger than that can also lead to faster restore times in many cases. Of course, values that are too high will lead to decreased performance because of thrashing.

Only the custom and directory archive formats are supported with this option. The input must be a regular file or directory (not, for example, a pipe). This option is ignored when emitting a script rather than connecting directly to a database server. Also, multiple jobs cannot be used together with the option `--single-transaction`.

-\l
--list

List the contents of the archive. The output of this operation can be used as input to the `-L` option. Note that if filtering switches such as `-n` or `-t` are used with `-l`, they will restrict the items listed.

-\L list-file
--use-list=list-file

Restore only those archive elements that are listed in `list-file`, and restore them in the order they appear in the file. Note that if filtering switches such as `-n` or `-t` are used with `-L`, they will further restrict the items restored.

`list-file` is normally created by editing the output of a previous `-l` operation. Lines can be moved or removed, and can also be commented out by placing a semicolon (`;`) at the start of the line. See below for examples.

-\n namespace
--schema=namespace

Restore only objects that are in the named schema. Multiple schemas may be specified with multiple `-n` switches. This can be combined with the `-t` option to restore just a specific table.

-\O
--no-owner

Do not output commands to set ownership of objects to match the original database. By default, `pg_restore` issues `ALTER OWNER` or `SET SESSION AUTHORIZATION` statements to set ownership of created schema elements. These statements will fail unless the initial connection to the database is made by a superuser (or the same user that owns all of the objects in the script). With `-O`, any user name can be used for the initial connection, and this user will own all the created objects.

-\P function-name(argtype [, ...])
--function=function-name(argtype [, ...])

Restore the named function only. Be careful to spell the function name and arguments exactly as they appear in the dump file’s table of contents. Multiple functions may be specified with multiple `-P` switches.

-\R
--no-reconnect

This option is obsolete but still accepted for backwards compatibility.

-\s
--schema-only

Restore only the schema (data definitions), not data, to the extent that schema entries are present in the archive.
This option is the inverse of --data-only. It is similar to, but for historical reasons not identical to, specifying --section=pre-data --section=post-data.

(Do not confuse this with the --schema option, which uses the word “schema” in a different meaning.)

-S username
--superuser=username

Specify the superuser user name to use when disabling triggers. This is relevant only if --disable-triggers is used.

-t table
--table=table

Restore definition and/or data of only the named table. For this purpose, “table” includes views, materialized views, sequences, and foreign tables. Multiple tables can be selected by writing multiple -t switches. This option can be combined with the -n option to specify table(s) in a particular schema.

Note: When -t is specified, pg_restore makes no attempt to restore any other database objects that the selected table(s) might depend upon. Therefore, there is no guarantee that a specific-table restore into a clean database will succeed.

Note: This flag does not behave identically to the -t flag of pg_dump. There is not currently any provision for wild-card matching in pg_restore, nor can you include a schema name within its -t.

Note: In versions prior to PostgreSQL 9.6, this flag matched only tables, not any other type of relation.

-T trigger
--trigger=trigger

Restore named trigger only. Multiple triggers may be specified with multiple -T switches.

-v
--verbose

Specifies verbose mode.

-V
--version

Print the pg_restore version and exit.

-x
--no-privileges
--no-acl

Prevent restoration of access privileges (grant/revoke commands).
--single-transaction

Execute the restore as a single transaction (that is, wrap the emitted commands in 
\texttt{BEGIN/COMMIT}). This ensures that either all the commands complete successfully, or no 
changes are applied. This option implies \texttt{--exit-on-error}.

--disable-triggers

This option is relevant only when performing a data-only restore. It instructs \texttt{pg_restore} to ex-
cute commands to temporarily disable triggers on the target tables while the data is reloaded. 
Use this if you have referential integrity checks or other triggers on the tables that you do not 
want to invoke during data reload.

Presently, the commands emitted for \texttt{--disable-triggers} must be done as superuser. So you 
should also specify a superuser name with \texttt{-S} or, preferably, run \texttt{pg_restore} as a PostgreSQL 
superuser.

--enable-row-security

This option is relevant only when restoring the contents of a table which has row security. By 
default, \texttt{pg_restore} will set row\_security to off, to ensure that all data is restored in to the table. If 
the user does not have sufficient privileges to bypass row security, then an error is thrown. This 
parameter instructs \texttt{pg_restore} to set row\_security to on instead, allowing the user to attempt to 
restore the contents of the table with row security enabled. This might still fail if the user does 
not have the right to insert the rows from the dump into the table.

Note that this option currently also requires the dump be in \texttt{INSERT} format, as \texttt{COPY FROM} does 
not support row security.

--if-exists

Use conditional commands (i.e. add an \texttt{IF EXISTS} clause) when cleaning database objects. This 
option is not valid unless \texttt{--clean} is also specified.

--no-data-for-failed-tables

By default, table data is restored even if the creation command for the table failed (e.g., because 
it already exists). With this option, data for such a table is skipped. This behavior is useful if 
the target database already contains the desired table contents. For example, auxiliary tables 
for PostgreSQL extensions such as PostGIS might already be loaded in the target database; 
specifying this option prevents duplicate or obsolete data from being loaded into them.

This option is effective only when restoring directly into a database, not when producing SQL 
script output.

--no-security-labels

Do not output commands to restore security labels, even if the archive contains them.

--no-tablespaces

Do not output commands to select tablespaces. With this option, all objects will be created in 
whichever tablespace is the default during restore.

--section=sectionname

Only restore the named section. The section name can be \texttt{pre-data}, \texttt{data}, or \texttt{post-data}. This 
option can be specified more than once to select multiple sections. The default is to restore all 
sections.
The data section contains actual table data as well as large-object definitions. Post-data items consist of definitions of indexes, triggers, rules and constraints other than validated check constraints. Pre-data items consist of all other data definition items.

--strict-names

Require that each schema (-n/--schema) and table (-t/--table) qualifier match at least one schema/table in the backup file.

--use-set-session-authorization

Output SQL-standard SET SESSION AUTHORIZATION commands instead of ALTER OWNER commands to determine object ownership. This makes the dump more standards-compatible, but depending on the history of the objects in the dump, might not restore properly.

-?
--help

Show help about pg_restore command line arguments, and exit.

pg_restore also accepts the following command line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket. The default is taken from the PGHOST environment variable, if set, else a Unix domain socket connection is attempted.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections. Defaults to the PGPORT environment variable, if set, or a compiled-in default.

-U username
--username=username

User name to connect as.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force pg_restore to prompt for a password before connecting to a database.

This option is never essential, since pg_restore will automatically prompt for a password if the server demands password authentication. However, pg_restore will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing -W to avoid the extra connection attempt.

--role=rolename

Specifies a role name to be used to perform the restore. This option causes pg_restore to issue a SET ROLE rolename command after connecting to the database. It is useful when the authenticated user (specified by -U) lacks privileges needed by pg_restore, but can switch to a role with
the required rights. Some installations have a policy against logging in directly as a superuser, and use of this option allows restores to be performed without violating the policy.

Environment

PGHOST
PGOPTIONS
PGPORT
PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14). However, it does not read PGDATABASE when a database name is not supplied.

Diagnostics

When a direct database connection is specified using the -d option, pg_restore internally executes SQL statements. If you have problems running pg_restore, make sure you are able to select information from the database using, for example, psql. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Notes

If your installation has any local additions to the template1 database, be careful to load the output of pg_restore into a truly empty database; otherwise you are likely to get errors due to duplicate definitions of the added objects. To make an empty database without any local additions, copy from template0 not template1, for example:

CREATE DATABASE foo WITH TEMPLATE template0;

The limitations of pg_restore are detailed below.

• When restoring data to a pre-existing table and the option --disable-triggers is used, pg_restore emits commands to disable triggers on user tables before inserting the data, then emits commands to re-enable them after the data has been inserted. If the restore is stopped in the middle, the system catalogs might be left in the wrong state.

• pg_restore cannot restore large objects selectively; for instance, only those for a specific table. If an archive contains large objects, then all large objects will be restored, or none of them if they are excluded via -L, -t, or other options.

See also the pg_dump documentation for details on limitations of pg_dump.

Once restored, it is wise to run ANALYZE on each restored table so the optimizer has useful statistics; see Section 24.1.3 and Section 24.1.6 for more information.
Examples

Assume we have dumped a database called `mydb` into a custom-format dump file:

```bash
$ pg_dump -Fc mydb > db.dump
```

To drop the database and recreate it from the dump:

```bash
$ dropdb mydb
$ pg_restore -C -d postgres db.dump
```

The database named in the `-d` switch can be any database existing in the cluster; `pg_restore` only uses it to issue the `CREATE DATABASE` command for `mydb`. With `-C`, data is always restored into the database name that appears in the dump file.

To reload the dump into a new database called `newdb`:

```bash
$ createdb -T template0 newdb
$ pg_restore -d newdb db.dump
```

Notice we don’t use `-C`, and instead connect directly to the database to be restored into. Also note that we clone the new database from `template0` not `template1`, to ensure it is initially empty.

To reorder database items, it is first necessary to dump the table of contents of the archive:

```bash
$ pg_restore -l db.dump > db.list
```

The listing file consists of a header and one line for each item, e.g.:

```
; Archive created at Mon Sep 14 13:55:39 2009
; dbname: DBDEMOS
; TOC Entries: 81
; Compression: 9
; Dump Version: 1.10-0
; Format: CUSTOM
; Integer: 4 bytes
; Offset: 8 bytes
; Dumped from database version: 8.3.5
; Dumped by pg_dump version: 8.3.8
;
; Selected TOC Entries:
;
3; 2615 2200 SCHEMA - public pasha
1861; 0 0 COMMENT - SCHEMA public pasha
1862; 0 0 ACL - public pasha
317; 1247 17715 TYPE public composite pasha
319; 1247 25899 DOMAIN public domain0 pasha
```

Semicolons start a comment, and the numbers at the start of lines refer to the internal archive ID assigned to each item.

Lines in the file can be commented out, deleted, and reordered. For example:

```
10; 145433 TABLE map_resolutions postgres
;2; 145344 TABLE species postgres
;4; 145359 TABLE nt_header postgres
```
could be used as input to `pg_restore` and would only restore items 10 and 6, in that order:

```
$ pg_restore -L db.list db.dump
```

See Also

`pg_dump`, `pg_dumpall`, `psql`
psql

Name

psql — PostgreSQL interactive terminal

Synopsis

psql [option...] [dbname [username]]

Description

psql is a terminal-based front-end to PostgreSQL. It enables you to type in queries interactively, issue them to PostgreSQL, and see the query results. Alternatively, input can be from a file or from command line arguments. In addition, psql provides a number of meta-commands and various shell-like features to facilitate writing scripts and automating a wide variety of tasks.

Options

-a
--echo-all

Print all nonempty input lines to standard output as they are read. (This does not apply to lines read interactively.) This is equivalent to setting the variable ECHO to all.

-A
--no-align

Switches to unaligned output mode. (The default output mode is otherwise aligned.)

-b
--echo-errors

Print failed SQL commands to standard error output. This is equivalent to setting the variable ECHO to errors.

-c command
--command=command

Specifies that psql is to execute the given command string, command. This option can be repeated and combined in any order with the -f option. When either -c or -f is specified, psql does not read commands from standard input; instead it terminates after processing all the -c and -f options in sequence.

command must be either a command string that is completely parsable by the server (i.e., it contains no psql-specific features), or a single backslash command. Thus you cannot mix SQL and psql meta-commands within a -c option. To achieve that, you could use repeated -c options or pipe the string into psql, for example:

psql -c '\x' -c 'SELECT * FROM foo;'
or

echo '\x \ SELECT * FROM foo;' | psql
\ is the separator meta-command.)

Each SQL command string passed to -c is sent to the server as a single query. Because of this, the server executes it as a single transaction even if the string contains multiple SQL commands, unless there are explicit BEGIN/COMMIT commands included in the string to divide it into multiple transactions. Also, psql only prints the result of the last SQL command in the string. This is different from the behavior when the same string is read from a file or fed to psql’s standard input, because then psql sends each SQL command separately.

Because of this behavior, putting more than one command in a single -c string often has unexpected results. It’s better to use repeated -c commands or feed multiple commands to psql’s standard input, either using echo as illustrated above, or via a shell here-document, for example:

```
psql <<EOF
\x
SELECT * FROM foo;
EOF
```

-d dbname
--dbname=dbname

Specifies the name of the database to connect to. This is equivalent to specifying dbname as the first non-option argument on the command line.

If this parameter contains an = sign or starts with a valid URI prefix (postgresql:// or postgres://), it is treated as a conninfo string. See Section 32.1.1 for more information.

-e
--echo-queries

Copy all SQL commands sent to the server to standard output as well. This is equivalent to setting the variable ECHO to queries.

-E
--echo-hidden

Echo the actual queries generated by \d and other backslash commands. You can use this to study psql’s internal operations. This is equivalent to setting the variable ECHO_HIDDEN to on.

-f filename
--file=filename

Read commands from the file filename, rather than standard input. This option can be repeated and combined in any order with the -c option. When either -c or -f is specified, psql does not read commands from standard input; instead it terminates after processing all the -c and -f options in sequence. Except for that, this option is largely equivalent to the meta-command !.

If filename is - (hyphen), then standard input is read until an EOF indication or \q meta-command. This can be used to intersperse interactive input with input from files. Note however that Readline is not used in this case (much as if -n had been specified).

Using this option is subtly different from writing psql < filename. In general, both will do what you expect, but using -f enables some nice features such as error messages with line numbers. There is also a slight chance that using this option will reduce the start-up overhead.

On the other hand, the variant using the shell’s input redirection is (in theory) guaranteed to yield exactly the same output you would have received had you entered everything by hand.

-F separator
--field-separator=separator

Use separator as the field separator for unaligned output. This is equivalent to \pset fieldsep or \f.
-h hostname
--host=hostname

Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix-domain socket.

-H
--html

Turn on HTML tabular output. This is equivalent to \psql format html or the \H command.

-l
--list

List all available databases, then exit. Other non-connection options are ignored. This is similar to the meta-command \list.

-L filename
--log-file=filename

Write all query output into file filename, in addition to the normal output destination.

-n
--no-readline

Do not use Readline for line editing and do not use the command history. This can be useful to turn off tab expansion when cutting and pasting.

-o filename
--output=filename

Put all query output into file filename. This is equivalent to the command \o.

-P port
--port=port

Specifies the TCP port or the local Unix-domain socket file extension on which the server is listening for connections. Defaults to the value of the PGPORT environment variable or, if not set, to the port specified at compile time, usually 5432.

-P assignment
--pset=assignment

Specifies printing options, in the style of \pset. Note that here you have to separate name and value with an equal sign instead of a space. For example, to set the output format to LaTeX, you could write -P format=latex.

-q
--quiet

Specifies that psql should do its work quietly. By default, it prints welcome messages and various informational output. If this option is used, none of this happens. This is useful with the -c option. This is equivalent to setting the variable QUIET to on.

-R separator
--record-separator=separator

Use separator as the record separator for unaligned output. This is equivalent to the \ps recordsep command.
-s
--single-step

Run in single-step mode. That means the user is prompted before each command is sent to the server, with the option to cancel execution as well. Use this to debug scripts.

-S
--single-line

Runs in single-line mode where a newline terminates an SQL command, as a semicolon does.

**Note:** This mode is provided for those who insist on it, but you are not necessarily encouraged to use it. In particular, if you mix SQL and meta-commands on a line the order of execution might not always be clear to the inexperienced user.

-t
--tuples-only

Turn off printing of column names and result row count footers, etc. This is equivalent to the \t command.

-T table_options
--table-attr=table_options

Specifies options to be placed within the HTML table tag. See \pset for details.

-U username
--username=username

Connect to the database as the user *username* instead of the default. (You must have permission to do so, of course.)

-v assignment
--set=assignment
--variable=assignment

Perform a variable assignment, like the \set meta-command. Note that you must separate name and value, if any, by an equal sign on the command line. To unset a variable, leave off the equal sign. To set a variable with an empty value, use the equal sign but leave off the value. These assignments are done during a very early stage of start-up, so variables reserved for internal purposes might get overwritten later.

-V
--version

Print the psql version and exit.

-w
--no-password

Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a .pgpass file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

Note that this option will remain set for the entire session, and so it affects uses of the meta-command \connect as well as the initial connection attempt.

-W
--password

Force psql to prompt for a password before connecting to a database.
This option is never essential, since psql will automatically prompt for a password if the server demands password authentication. However, psql will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing \W to avoid the extra connection attempt.

Note that this option will remain set for the entire session, and so it affects uses of the meta-command \connect as well as the initial connection attempt.

-x
--expanded

Turn on the expanded table formatting mode. This is equivalent to the \x command.

-X, --no-psqlrc

Do not read the start-up file (neither the system-wide psqlrc file nor the user’s ~/.psqlrc file).

-z
--field-separator-zero

Set the field separator for unaligned output to a zero byte.

-0
--record-separator-zero

Set the record separator for unaligned output to a zero byte. This is useful for interfacing, for example, with xargs -0.

-1
--single-transaction

This option can only be used in combination with one or more -c and/or -f options. It causes psql to issue a BEGIN command before the first such option and a COMMIT command after the last one, thereby wrapping all the commands into a single transaction. This ensures that either all the commands complete successfully, or no changes are applied.

If the commands themselves contain BEGIN, COMMIT, or ROLLBACK, this option will not have the desired effects. Also, if an individual command cannot be executed inside a transaction block, specifying this option will cause the whole transaction to fail.

-?
--help[=topic]

Show help about psql and exit. The optional topic parameter (defaulting to options) selects which part of psql is explained: commands describes psql’s backslash commands; options describes the command-line options that can be passed to psql; and variables shows help about psql configuration variables.

Exit Status

psql returns 0 to the shell if it finished normally, 1 if a fatal error of its own occurs (e.g. out of memory, file not found), 2 if the connection to the server went bad and the session was not interactive, and 3 if an error occurred in a script and the variable ON_ERROR_STOP was set.
Usage

Connecting to a Database

psql is a regular PostgreSQL client application. In order to connect to a database you need to know the name of your target database, the host name and port number of the server, and what user name you want to connect as. psql can be told about those parameters via command line options, namely \texttt{-d}, \texttt{-h}, \texttt{-p}, and \texttt{-U} respectively. If an argument is found that does not belong to any option it will be interpreted as the database name (or the user name, if the database name is already given). Not all of these options are required; there are useful defaults. If you omit the host name, psql will connect via a Unix-domain socket to a server on the local host, or via TCP/IP to \texttt{localhost} on machines that don't have Unix-domain sockets. The default port number is determined at compile time. Since the database server uses the same default, you will not have to specify the port in most cases. The default user name is your operating-system user name, as is the default database name. Note that you cannot just connect to any database under any user name. Your database administrator should have informed you about your access rights.

When the defaults aren't quite right, you can save yourself some typing by setting the environment variables \texttt{PGDATABASE}, \texttt{PGHOST}, \texttt{PGPORT} and/or \texttt{PGUSER} to appropriate values. (For additional environment variables, see Section 32.14.) It is also convenient to have a \texttt{~/.pgpass} file to avoid regularly having to type in passwords. See Section 32.15 for more information.

An alternative way to specify connection parameters is in a \texttt{conninfo} string or a URI, which is used instead of a database name. This mechanism give you very wide control over the connection. For example:

\begin{verbatim}
$ psql "service=myservice sslmode=require"
$ psql postgresql://dbmaster:5433/mydb?sslmode=require
\end{verbatim}

This way you can also use LDAP for connection parameter lookup as described in Section 32.17. See Section 32.1.2 for more information on all the available connection options.

If the connection could not be made for any reason (e.g., insufficient privileges, server is not running on the targeted host, etc.), psql will return an error and terminate.

If both standard input and standard output are a terminal, then psql sets the client encoding to “auto”, which will detect the appropriate client encoding from the locale settings (\texttt{LC_CTYPE} environment variable on Unix systems). If this doesn’t work out as expected, the client encoding can be overridden using the environment variable \texttt{PGCLIENTENCODING}.

Entering SQL Commands

In normal operation, psql provides a prompt with the name of the database to which psql is currently connected, followed by the string \texttt{=>}. For example:

\begin{verbatim}
$ psql testdb
psql (9.6.17)
Type \texttt{*help*} for help.

testdb=>
\end{verbatim}

At the prompt, the user can type in SQL commands. Ordinarily, input lines are sent to the server when a command-terminating semicolon is reached. An end of line does not terminate a command. Thus
commands can be spread over several lines for clarity. If the command was sent and executed without error, the results of the command are displayed on the screen.

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin your session by removing publicly-writable schemas from `search_path`. One can add `options=-c search_path=\` to the connection string or issue `SELECT pg_catalog.set_config('search_path', ", false)` before other SQL commands. This consideration is not specific to `psql`; it applies to every interface for executing arbitrary SQL commands.

Whenever a command is executed, `psql` also polls for asynchronous notification events generated by `LISTEN` and `NOTIFY`.

While C-style block comments are passed to the server for processing and removal, SQL-standard comments are removed by `psql`.

**Meta-Commands**

 Anything you enter in `psql` that begins with an unquoted backslash is a `psql` meta-command that is processed by `psql` itself. These commands make `psql` more useful for administration or scripting. Meta-commands are often called slash or backslash commands.

 The format of a `psql` command is the backslash, followed immediately by a command verb, then any arguments. The arguments are separated from the command verb and each other by any number of whitespace characters.

 To include whitespace in an argument you can quote it with single quotes. To include a single quote in an argument, write two single quotes within single-quoted text. Anything contained in single quotes is furthermore subject to C-like substitutions for `\n` (new line), `\t` (tab), `\b` (backspace), `\r` (carriage return), `\f` (form feed), `\digs` (octal), and `\x` digits (hexadecimal). A backslash preceding any other character within single-quoted text quotes that single character, whatever it is.

 Within an argument, text that is enclosed in backquotes (``) is taken as a command line that is passed to the shell. The output of the command (with any trailing newline removed) replaces the backquoted text.

 If an unquoted colon (:) followed by a `psql` variable name appears within an argument, it is replaced by the variable’s value, as described in SQL Interpolation.

 Some commands take an SQL identifier (such as a table name) as argument. These arguments follow the syntax rules of SQL: Unquoted letters are forced to lowercase, while double quotes (" ) protect letters from case conversion and allow incorporation of whitespace into the identifier. Within double quotes, paired double quotes reduce to a single double quote in the resulting name. For example, `FOO"BAR"BAZ` is interpreted as `fooBARbaz`, and "A weird"" name" becomes `A weird" name`.

 Parsing for arguments stops at the end of the line, or when another unquoted backslash is found. An unquoted backslash is taken as the beginning of a new meta-command. The special sequence `\` (two backslashes) marks the end of arguments and continues parsing SQL commands, if any. That way SQL and `psql` commands can be freely mixed on a line. But in any case, the arguments of a meta-command cannot continue beyond the end of the line.

 The following meta-commands are defined:

 `\a`

 If the current table output format is unaligned, it is switched to aligned. If it is not unaligned, it is set to unaligned. This command is kept for backwards compatibility. See `\pset` for a more general solution.
Establishes a new connection to a PostgreSQL server. The connection parameters to use can be specified either using a positional syntax, or using conninfo connection strings as detailed in Section 32.1.1.

Where the command omits database name, user, host, or port, the new connection can reuse values from the previous connection. By default, values from the previous connection are reused except when processing a conninfo string. Passing a first argument of -reuse-previous=on or -reuse-previous=off overrides that default. When the command neither specifies nor reuses a particular parameter, the libpq default is used. Specifying any of dbname, username, host or port as - is equivalent to omitting that parameter.

If the new connection is successfully made, the previous connection is closed. If the connection attempt failed (wrong user name, access denied, etc.), the previous connection will only be kept if psql is in interactive mode. When executing a non-interactive script, processing will immediately stop with an error. This distinction was chosen as a user convenience against typos on the one hand, and a safety mechanism that scripts are not accidentally acting on the wrong database on the other hand.

Examples:

```
=> \c mydb myuser host.dom 6432
=> \c service=foo
=> \c "host=localhost port=5432 dbname=mydb connect_timeout=10 sslmode=disable"
=> \c postgresql://tom@localhost/mydb?application_name=myapp
```

\C [ title ]

Sets the title of any tables being printed as the result of a query or unset any such title. This command is equivalent to \pset title title. (The name of this command derives from "caption", as it was previously only used to set the caption in an HTML table.)

\cd [ directory ]

Changes the current working directory to directory. Without argument, changes to the current user’s home directory.

**Tip:** To print your current working directory, use \! pwd.

\conninfo

Outputs information about the current database connection.

\copy { table [ ( column_list ) ] | ( query ) } { from | to } { ’filename’ | program ’command’ | stdin | stdout | pstdin | pstdout } [ [ with ] ( option [, ...] ) ]

Performs a frontend (client) copy. This is an operation that runs an SQL COPY command, but instead of the server reading or writing the specified file, psql reads or writes the file and routes the data between the server and the local file system. This means that file accessibility and privileges are those of the local user, not the server, and no SQL superuser privileges are required.

When program is specified, command is executed by psql and the data passed from or to command is routed between the server and the client. Again, the execution privileges are those of the local user, not the server, and no SQL superuser privileges are required.
For `\copy` command, data rows are read from the same source that issued the command, continuing until `\.` is read or the stream reaches EOF. This option is useful for populating tables in-line within a SQL script file. For `\copy` command, output is sent to the same place as psql command output, and the `COPY count` command status is not printed (since it might be confused with a data row). To read/write psql’s standard input or output regardless of the current command source or `\o` option, write `from pstdin` or `to pstdout`.

The syntax of this command is similar to that of the SQL COPY command. All options other than the data source/destination are as specified for COPY. Because of this, special parsing rules apply to the `\copy` command. In particular, psql’s variable substitution rules and backslash escapes do not apply.

**Tip:** Another way to obtain the same result as `\copy` command and terminate it with `\g` or `\g |`. Unlike `\copy`, this method allows the command to span multiple lines; also, variable interpolation and backquote expansion can be used.

**Tip:** These operations are not as efficient as the SQL `COPY` command with a file or program data source or destination, because all data must pass through the client/server connection. For large amounts of data the SQL command might be preferable.

```
\copyright
```

Shows the copyright and distribution terms of PostgreSQL.

```
\crosstabview [ colV] [ colH [ colD [ sortcolH ] ] ]
```

Executes the current query buffer (like `\g`) and shows the results in a crosstab grid. The query must return at least three columns. The output column identified by `colV` becomes a vertical header and the output column identified by `colH` becomes a horizontal header. `colD` identifies the output column to display within the grid. `sortcolH` identifies an optional sort column for the horizontal header.

Each column specification can be a column number (starting at 1) or a column name. The usual SQL case folding and quoting rules apply to column names. If omitted, `colV` is taken as column 1 and `colH` as column 2. `colH` must differ from `colV` and it must differ from `colD`. If `colD` is not specified, then there must be exactly three columns in the query result, and the column that is neither `colV` nor `colH` is taken to be `colD`.

The vertical header, displayed as the leftmost column, contains the values found in column `colV`, in the same order as in the query results, but with duplicates removed.

The horizontal header, displayed as the first row, contains the values found in column `colH`, with duplicates removed. By default, these appear in the same order as in the query results. But if the optional `sortcolH` argument is given, it identifies a column whose values must be integer numbers, and the values from `colH` will appear in the horizontal header sorted according to the corresponding `sortcolH` values.

Inside the crosstab grid, for each distinct value `x` of `colH` and each distinct value `y` of `colV`, the cell located at the intersection `(x, y)` contains the value of the `colD` column in the query result row for which the value of `colH` is `x` and the value of `colV` is `y`. If there is no such row, the cell is empty. If there are multiple such rows, an error is reported.
The command form \d+ is identical, except that more information is displayed: any comments associated with the columns of the table are shown, as is the presence of OIDs in the table, the view definition if the relation is a view, a non-default replica identity setting.

By default, only user-created objects are shown; supply a pattern or the $ modifier to include system objects.

**Note:** If \d is used without a pattern argument, it is equivalent to \dtvmsE which will show a list of all visible tables, views, materialized views, sequences and foreign tables. This is purely a convenience measure.

\da[S] [ pattern ]

Lists aggregate functions, together with their return type and the data types they operate on. If pattern is specified, only aggregates whose names match the pattern are shown. By default, only user-created objects are shown; supply a pattern or the $ modifier to include system objects.

\dA[+] [ pattern ]

Lists access methods. If pattern is specified, only access methods whose names match the pattern are shown. If + is appended to the command name, each access method is listed with its associated handler function and description.

\db[+] [ pattern ]

Lists tablespaces. If pattern is specified, only tablespaces whose names match the pattern are shown. If + is appended to the command name, each tablespace is listed with its associated options, on-disk size, permissions and description.

\dc[S+] [ pattern ]

Lists conversions between character-set encodings. If pattern is specified, only conversions whose names match the pattern are listed. By default, only user-created objects are shown; supply a pattern or the $ modifier to include system objects. If + is appended to the command name, each object is listed with its associated description.

\dC[+] [ pattern ]

Lists type casts. If pattern is specified, only casts whose source or target types match the pattern are listed. If + is appended to the command name, each object is listed with its associated description.

\dd[S] [ pattern ]

Shows the descriptions of objects of type constraint, operator class, operator family, rule, and trigger. All other comments may be viewed by the respective backslash commands for those object types.
\dd displays descriptions for objects matching the \texttt{pattern}, or of visible objects of the appropriate type if no argument is given. But in either case, only objects that have a description are listed. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects.

Descriptions for objects can be created with the \texttt{COMMENT} SQL command.

\texttt{\ddp [ \textit{pattern} ]}

Lists default access privilege settings. An entry is shown for each role (and schema, if applicable) for which the default privilege settings have been changed from the built-in defaults. If \textit{pattern} is specified, only entries whose role name or schema name matches the pattern are listed.

The ALTER DEFAULT PRIVILEGES command is used to set default access privileges. The meaning of the privilege display is explained under GRANT.

\texttt{\dD[S+] [ \textit{pattern} ]}

Lists domains. If \textit{pattern} is specified, only domains whose names match the pattern are shown. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects. If + is appended to the command name, each object is listed with its associated permissions and description.

\texttt{\dE[S+] [ \textit{pattern} ]}
\texttt{\di[S+] [ \textit{pattern} ]}
\texttt{\dm[S+] [ \textit{pattern} ]}
\texttt{\ds[S+] [ \textit{pattern} ]}
\texttt{\dt[S+] [ \textit{pattern} ]}
\texttt{\dv[S+] [ \textit{pattern} ]}

In this group of commands, the letters \texttt{E}, \texttt{i}, \texttt{m}, \texttt{s}, \texttt{t}, and \texttt{v} stand for foreign table, index, materialized view, sequence, table, and view, respectively. You can specify any or all of these letters, in any order, to obtain a listing of objects of these types. For example, \texttt{\dit} lists indexes and tables. If + is appended to the command name, each object is listed with its physical size on disk and its associated description, if any. If \textit{pattern} is specified, only objects whose names match the pattern are listed. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects.

\texttt{\des[+] [ \textit{pattern} ]}

Lists foreign servers (mnemonic: “external servers”). If \textit{pattern} is specified, only those servers whose name matches the pattern are listed. If the form \texttt{\des+} is used, a full description of each server is shown, including the server’s ACL, type, version, options, and description.

\texttt{\det[+] [ \textit{pattern} ]}

Lists foreign tables (mnemonic: “external tables”). If \textit{pattern} is specified, only entries whose table name or schema name matches the pattern are listed. If the form \texttt{\det+} is used, generic options and the foreign table description are also displayed.

\texttt{\deu[+] [ \textit{pattern} ]}

Lists user mappings (mnemonic: “external users”). If \textit{pattern} is specified, only those mappings whose user names match the pattern are listed. If the form \texttt{\deu+} is used, additional information about each mapping is shown.

\textbf{Caution}

\texttt{\deu+} might also display the user name and password of the remote user, so care should be taken not to disclose them.
\texttt{\textbackslash dew[+] [ pattern ]}

Lists foreign-data wrappers (mnemonic: “external wrappers”). If \texttt{pattern} is specified, only those foreign-data wrappers whose name matches the pattern are listed. If the form \texttt{\textbackslash dew+} is used, the ACL, options, and description of the foreign-data wrapper are also shown.

\texttt{\textbackslash df[antwS+] [ pattern ]}

Lists functions, together with their result data types, argument data types, and function types, which are classified as “agg” (aggregate), “normal”, “trigger”, or “window”. To display only functions of specific type(s), add the corresponding letters \texttt{a}, \texttt{n}, \texttt{t}, or \texttt{w} to the command. If \texttt{pattern} is specified, only functions whose names match the pattern are shown. By default, only user-created objects are shown; supply a pattern or the \texttt{S} modifier to include system objects. If the form \texttt{\textbackslash df+} is used, additional information about each function is shown, including volatility, parallel safety, owner, security classification, access privileges, language, source code and description.

\textbf{Tip:} To look up functions taking arguments or returning values of a specific data type, use your pager’s search capability to scroll through the \texttt{\textbackslash df} output.

\texttt{\textbackslash dF[+] [ pattern ]}

Lists text search configurations. If \texttt{pattern} is specified, only configurations whose names match the pattern are shown. If the form \texttt{\textbackslash dF+} is used, a full description of each configuration is shown, including the underlying text search parser and the dictionary list for each parser token type.

\texttt{\textbackslash dFd[+] [ pattern ]}

Lists text search dictionaries. If \texttt{pattern} is specified, only dictionaries whose names match the pattern are shown. If the form \texttt{\textbackslash dFd+} is used, additional information is shown about each selected dictionary, including the underlying text search template and the option values.

\texttt{\textbackslash dFn[+] [ pattern ]}

Lists text search parsers. If \texttt{pattern} is specified, only parsers whose names match the pattern are shown. If the form \texttt{\textbackslash dFn+} is used, a full description of each parser is shown, including the underlying functions and the list of recognized token types.

\texttt{\textbackslash dFt[+] [ pattern ]}

Lists text search templates. If \texttt{pattern} is specified, only templates whose names match the pattern are shown. If the form \texttt{\textbackslash dFt+} is used, additional information is shown about each template, including the underlying function names.

\texttt{\textbackslash dg[S+] [ pattern ]}

Lists database roles. (Since the concepts of “users” and “groups” have been unified into “roles”, this command is now equivalent to \texttt{\textbackslash du}.) By default, only user-created roles are shown; supply the \texttt{S} modifier to include system roles. If \texttt{pattern} is specified, only those roles whose names match the pattern are listed. If the form \texttt{\textbackslash dg+} is used, additional information is shown about each role; currently this adds the comment for each role.

\texttt{\textbackslash dl}

This is an alias for \texttt{\textbackslash lo_list}, which shows a list of large objects.

\texttt{\textbackslash dL[S+] [ pattern ]}

Lists procedural languages. If \texttt{pattern} is specified, only languages whose names match the pattern are listed. By default, only user-created languages are shown; supply the \texttt{S} modifier to
include system objects. If + is appended to the command name, each language is listed with its
call handler, validator, access privileges, and whether it is a system object.

\dn [S+] [ pattern ]

Lists schemas (namespaces). If pattern is specified, only schemas whose names match the
pattern are listed. By default, only user-created objects are shown; supply a pattern or the S
modifier to include system objects. If + is appended to the command name, each object is listed
with its associated permissions and description, if any.

\do [S+] [ pattern ]

Lists operators with their operand and result types. If pattern is specified, only operators whose
names match the pattern are listed. By default, only user-created objects are shown; supply a
pattern or the S modifier to include system objects. If + is appended to the command name,
additional information about each operator is shown, currently just the name of the underlying
function.

\dO [S+] [ pattern ]

Lists collations. If pattern is specified, only collations whose names match the pattern are
listed. By default, only user-created objects are shown; supply a pattern or the S modifier to
include system objects. If + is appended to the command name, each collation is listed with
its associated description, if any. Note that only collations usable with the current database’s
encoding are shown, so the results may vary in different databases of the same installation.

\dp [ pattern ]

Lists tables, views and sequences with their associated access privileges. If pattern is specified,
only tables, views and sequences whose names match the pattern are listed.

The GRANT and REVOKE commands are used to set access privileges. The meaning of the
privilege display is explained under GRANT.

\drds [ role-pattern [ database-pattern ] ]

Lists defined configuration settings. These settings can be role-specific, database-specific, or
both. role-pattern and database-pattern are used to select specific roles and databases
to list, respectively. If omitted, or if * is specified, all settings are listed, including those not
role-specific or database-specific, respectively.

The ALTER ROLE and ALTER DATABASE commands are used to define per-role and per-
database configuration settings.

\dT [S+] [ pattern ]

Lists data types. If pattern is specified, only types whose names match the pattern are listed.
If + is appended to the command name, each type is listed with its internal name and size, its
allowed values if it is an enum type, and its associated permissions. By default, only user-created
objects are shown; supply a pattern or the S modifier to include system objects.

\du [S+] [ pattern ]

Lists database roles. (Since the concepts of “users” and “groups” have been unified into “roles”,
this command is now equivalent to \dg.) By default, only user-created roles are shown; supply
the S modifier to include system roles. If pattern is specified, only those roles whose names
match the pattern are listed. If the form \du+ is used, additional information is shown about eachole; currently this adds the comment for each role.

\dx [+] [ pattern ]

Lists installed extensions. If pattern is specified, only those extensions whose names match the
pattern are listed. If the form \dx+ is used, all the objects belonging to each matching extension
are listed.

\dy{[+] [ pattern ]}

Lists event triggers. If \texttt{pattern} is specified, only those event triggers whose names match the pattern are listed. If + is appended to the command name, each object is listed with its associated description.

\e or \edit [ filename ] [ line_number ]

If \texttt{filename} is specified, the file is edited; after the editor exits, its content is copied back to the query buffer. If no \texttt{filename} is given, the current query buffer is copied to a temporary file which is then edited in the same fashion.

The new query buffer is then re-parsed according to the normal rules of \texttt{psql}, where the whole buffer is treated as a single line. (Thus you cannot make scripts this way. Use \texttt{\i} for that.) This means that if the query ends with (or contains) a semicolon, it is immediately executed. Otherwise it will merely wait in the query buffer; type semicolon or \texttt{\g} to send it, or \texttt{\r} to cancel.

If a line number is specified, \texttt{psql} will position the cursor on the specified line of the file or query buffer. Note that if a single all-digits argument is given, \texttt{psql} assumes it is a line number, not a file name.

\textbf{Tip:} See under \textit{Environment} for how to configure and customize your editor.

\echo text [ ... ]

Prints the arguments to the standard output, separated by one space and followed by a newline. This can be useful to intersperse information in the output of scripts. For example:

\begin{verbatim}
 => \echo 'date'
 Tue Oct 26 21:40:57 CEST 1999
\end{verbatim}

If the first argument is an unquoted \texttt{-n} the trailing newline is not written.

\textbf{Tip:} If you use the \texttt{\o} command to redirect your query output you might wish to use \texttt{\qecho} instead of this command.

\ef [ function_description [ line_number ] ]

This command fetches and edits the definition of the named function, in the form of a \texttt{CREATE OR REPLACE FUNCTION} command. Editing is done in the same way as for \texttt{\edit}. After the editor exits, the updated command waits in the query buffer; type semicolon or \texttt{\g} to send it, or \texttt{\r} to cancel.

The target function can be specified by name alone, or by name and arguments, for example \texttt{foo(integer, text)}. The argument types must be given if there is more than one function of the same name.

If no function is specified, a blank \texttt{CREATE FUNCTION} template is presented for editing.

If a line number is specified, \texttt{psql} will position the cursor on the specified line of the function body. (Note that the function body typically does not begin on the first line of the file.)

\textbf{Tip:} See under \textit{Environment} for how to configure and customize your editor.
\encoding [ encoding ]

Sets the client character set encoding. Without an argument, this command shows the current encoding.

\errverbose

Repeats the most recent server error message at maximum verbosity, as though \texttt{VERBOSITY} were set to \texttt{verbose} and \texttt{SHOW\_CONTEXT} were set to \texttt{always}.

\ev [ view_name [ line_number ] ]

This command fetches and edits the definition of the named view, in the form of a \texttt{CREATE OR REPLACE VIEW} command. Editing is done in the same way as for \texttt{\edit}. After the editor exits, the updated command waits in the query buffer; type \texttt{\g} to send it, or \texttt{\r} to cancel.

If no view is specified, a blank \texttt{CREATE VIEW} template is presented for editing.

If a line number is specified, psql will position the cursor on the specified line of the view definition.

\f [ string ]

Sets the field separator for unaligned query output. The default is the vertical bar (|). See also \texttt{\pset} for a generic way of setting output options.

\g [ filename ]
\g [ command ]

Sends the current query input buffer to the server, and optionally stores the query’s output in \texttt{filename} or pipes the output to the shell command \texttt{command}. The file or command is written to only if the query successfully returns zero or more tuples, not if the query fails or is a non-data-returning SQL command.

A bare \texttt{\g} is essentially equivalent to a semicolon. A \texttt{\g} with argument is a “one-shot” alternative to the \texttt{\o} command.

\gexec

Sends the current query input buffer to the server, then treats each column of each row of the query’s output (if any) as a SQL statement to be executed. For example, to create an index on each column of \texttt{my_table}:

\begin{verbatim}
  => SELECT format('create index on my_table(%I)', attname)
  => FROM pg_attribute
  => WHERE attrelid = 'my_table'::regclass AND attnum > 0
  => ORDER BY attnum
  => \gexec
  CREATE INDEX
  CREATE INDEX
  CREATE INDEX
  CREATE INDEX
\end{verbatim}

The generated queries are executed in the order in which the rows are returned, and left-to-right within each row if there is more than one column. NULL fields are ignored. The generated queries are sent literally to the server for processing, so they cannot be psql meta-commands nor contain psql variable references. If any individual query fails, execution of the remaining queries continues unless \texttt{ON\_ERROR\_STOP} is set. Execution of each query is subject to \texttt{ECHO} processing. (Setting \texttt{ECHO} to \texttt{all} or \texttt{queries} is often advisable when using \texttt{\gexec}.) Query logging, single-step mode, timing, and other query execution features apply to each generated query as well.
\gset [ prefix ]

Sends the current query input buffer to the server and stores the query’s output into psql variables (see Variables). The query to be executed must return exactly one row. Each column of the row is stored into a separate variable, named the same as the column. For example:

```
=> SELECT 'hello' AS var1, 10 AS var2
=> \gset
=> \echo :var1 :var2
hello 10
```

If you specify a prefix, that string is prepended to the query’s column names to create the variable names to use:

```
=> SELECT 'hello' AS var1, 10 AS var2
=> \gset result_
=> \echo :result_var1 :result_var2
hello 10
```

If a column result is NULL, the corresponding variable is unset rather than being set.

If the query fails or does not return one row, no variables are changed.

\h or \help [ command ]

Gives syntax help on the specified SQL command. If command is not specified, then psql will list all the commands for which syntax help is available. If command is an asterisk (*), then syntax help on all SQL commands is shown.

**Note:** To simplify typing, commands that consists of several words do not have to be quoted. Thus it is fine to type \help alter table.

\H or \html

Turns on HTML query output format. If the HTML format is already on, it is switched back to the default aligned text format. This command is for compatibility and convenience, but see \pset about setting other output options.

\i or \include filename

Reads input from the file filename and executes it as though it had been typed on the keyboard. If filename is -, then standard input is read until an EOF indication or \q meta-command. This can be used to intersperse interactive input with input from files. Note that Readline behavior will be used only if it is active at the outermost level.

**Note:** If you want to see the lines on the screen as they are read you must set the variable ECHO to all.

\ir or \include_relative filename

The \ir command is similar to \i, but resolves relative file names differently. When executing in interactive mode, the two commands behave identically. However, when invoked from a script, \ir interprets file names relative to the directory in which the script is located, rather than the current working directory.
\l[+] or \list[+] [ pattern ]

List the databases in the server and show their names, owners, character set encodings, and access
privileges. If pattern is specified, only databases whose names match the pattern are listed. If +
is appended to the command name, database sizes, default tablespaces, and descriptions are also
displayed. (Size information is only available for databases that the current user can connect to.)

\lo_export oid filename

Reads the large object with OID oid from the database and writes it to filename. Note that
this is subtly different from the server function lo_export, which acts with the permissions of
the user that the database server runs as and on the server’s file system.

Tip: Use \lo_list to find out the large object’s OID.

\lo_import filename [ comment ]

Stores the file into a PostgreSQL large object. Optionally, it associates the given comment with
the object. Example:

foo=> \lo_import '/home/peter/pictures/photo.xcf' 'a picture of me'
lo_import 152801

The response indicates that the large object received object ID 152801, which can be used to
access the newly-created large object in the future. For the sake of readability, it is recommended
to always associate a human-readable comment with every object. Both OIDs and comments can
be viewed with the \lo_list command.

Note that this command is subtly different from the server-side lo_import because it acts as
the local user on the local file system, rather than the server’s user and file system.

\lo_list

Shows a list of all PostgreSQL large objects currently stored in the database, along with any
comments provided for them.

\lo_unlink oid

Deletes the large object with OID oid from the database.

Tip: Use \lo_list to find out the large object’s OID.

\o or \out [ filename ]
\o or \out [ command ]

Arranges to save future query results to the file filename or pipe future results to the shell
command command. If no argument is specified, the query output is reset to the standard output.
“Query results” includes all tables, command responses, and notices obtained from the database
server, as well as output of various backslash commands that query the database (such as \d),
but not error messages.

Tip: To intersperse text output in between query results, use \qecho.

\p or \print

Print the current query buffer to the standard output.
Changes the password of the specified user (by default, the current user). This command prompts for the new password, encrypts it, and sends it to the server as an ALTER ROLE command. This makes sure that the new password does not appear in cleartext in the command history, the server log, or elsewhere.

Prompts the user to supply text, which is assigned to the variable name. An optional prompt string, text, can be specified. (For multiword prompts, surround the text with single quotes.) By default, \prompt uses the terminal for input and output. However, if the -f command line switch was used, \prompt uses standard input and standard output.

This command sets options affecting the output of query result tables. option indicates which option is to be set. The semantics of value vary depending on the selected option. For some options, omitting value causes the option to be toggled or unset, as described under the particular option. If no such behavior is mentioned, then omitting value just results in the current setting being displayed. \pset without any arguments displays the current status of all printing options.

Adjustable printing options are:

border

The value must be a number. In general, the higher the number the more borders and lines the tables will have, but details depend on the particular format. In HTML format, this will translate directly into the border=... attribute. In most other formats only values 0 (no border), 1 (internal dividing lines), and 2 (table frame) make sense, and values above 2 will be treated the same as border = 2. The latex and latex-longtable formats additionally allow a value of 3 to add dividing lines between data rows.

columns

Sets the target width for the wrapped format, and also the width limit for determining whether output is wide enough to require the pager or switch to the vertical display in expanded auto mode. Zero (the default) causes the target width to be controlled by the environment variable COLUMNS, or the detected screen width if COLUMNS is not set. In addition, if columns is zero then the wrapped format only affects screen output. If columns is nonzero then file and pipe output is wrapped to that width as well.

expanded (or x)

If value is specified it must be either on or off, which will enable or disable expanded mode, or auto. If value is omitted the command toggles between the on and off settings. When expanded mode is enabled, query results are displayed in two columns, with the column name on the left and the data on the right. This mode is useful if the data wouldn’t fit on the screen in the normal “horizontal” mode. In the auto setting, the expanded mode is used whenever the query output has more than one column and is wider than the screen; otherwise, the regular mode is used. The auto setting is only effective in the aligned and wrapped formats. In other formats, it always behaves as if the expanded mode is off.

fieldsep

Specifies the field separator to be used in unaligned output format. That way one can create, for example, tab- or comma-separated output, which other programs might prefer. To set a
tab as field separator, type \pset fieldsep ' \t '. The default field separator is ' | ' (a vertical bar).

fieldsep_zero

Sets the field separator to use in unaligned output format to a zero byte.

footer

If value is specified it must be either on or off which will enable or disable display of the table footer (the (n rows) count). If value is omitted the command toggles footer display on or off.

format

Sets the output format to one of unaligned, aligned, wrapped, html, asciidoc, latex (uses tabular), latex-longtable, or troff-ms. Unique abbreviations are allowed.

unaligned format writes all columns of a row on one line, separated by the currently active field separator. This is useful for creating output that might be intended to be read in by other programs (for example, tab-separated or comma-separated format).

aligned format is the standard, human-readable, nicely formatted text output; this is the default.

wrapped format is like aligned but wraps wide data values across lines to make the output fit in the target column width. The target width is determined as described under the columns option. Note that psql will not attempt to wrap column header titles; therefore, wrapped format behaves the same as aligned if the total width needed for column headers exceeds the target.

The html, asciidoc, latex, latex-longtable, and troff-ms formats put out tables that are intended to be included in documents using the respective mark-up language. They are not complete documents! This might not be necessary in HTML, but in LaTeX you must have a complete document wrapper. latex-longtable also requires the LaTeX longtable and booktabs packages.

linestyle

Sets the border line drawing style to one of ascii, old-ascii, or unicode. Unique abbreviations are allowed. (That would mean one letter is enough.) The default setting is ascii. This option only affects the aligned and wrapped output formats.

ascii style uses plain ASCII characters. Newlines in data are shown using a + symbol in the right-hand margin. When the wrapped format wraps data from one line to the next without a newline character, a dot (.) is shown in the right-hand margin of the first line, and again in the left-hand margin of the following line.

old-ascii style uses plain ASCII characters, using the formatting style used in PostgreSQL 8.4 and earlier. Newlines in data are shown using a : symbol in place of the left-hand column separator. When the data is wrapped from one line to the next without a newline character, a ; symbol is used in place of the left-hand column separator.

unicode style uses Unicode box-drawing characters. Newlines in data are shown using a carriage return symbol in the right-hand margin. When the data is wrapped from one line to the next without a newline character, an ellipsis symbol is shown in the right-hand margin of the first line, and again in the left-hand margin of the following line.

When the border setting is greater than zero, the linestyle option also determines the characters with which the border lines are drawn. Plain ASCII characters work everywhere, but Unicode characters look nicer on displays that recognize them.
null

Sets the string to be printed in place of a null value. The default is to print nothing, which can easily be mistaken for an empty string. For example, one might prefer \pset null '(null)'.

numericlocale

If value is specified it must be either on or off which will enable or disable display of a locale-specific character to separate groups of digits to the left of the decimal marker. If value is omitted the command toggles between regular and locale-specific numeric output.

pager

Controls use of a pager program for query and psql help output. If the environment variable PAGER is set, the output is piped to the specified program. Otherwise a platform-dependent default (such as more) is used.

When the pager option is off, the pager program is not used. When the pager option is on, the pager is used when appropriate, i.e., when the output is to a terminal and will not fit on the screen. The pager option can also be set to always, which causes the pager to be used for all terminal output regardless of whether it fits on the screen. \pset pager without a value toggles pager use on and off.

pager_min_lines

If pager_min_lines is set to a number greater than the page height, the pager program will not be called unless there are at least this many lines of output to show. The default setting is 0.

recordsep

Specifies the record (line) separator to use in unaligned output format. The default is a newline character.

recordsep_zero

Sets the record separator to use in unaligned output format to a zero byte.

tableattr (or T)

In HTML format, this specifies attributes to be placed inside the table tag. This could for example be cellpadding or bgcolor. Note that you probably don’t want to specify border here, as that is already taken care of by \pset border. If no value is given, the table attributes are unset.

In latex-longtable format, this controls the proportional width of each column containing a left-aligned data type. It is specified as a whitespace-separated list of values, e.g. ‘0.2 0.2 0.6’. Unspecified output columns use the last specified value.

title (or C)

Sets the table title for any subsequently printed tables. This can be used to give your output descriptive tags. If no value is given, the title is unset.

tuples_only (or t)

If value is specified it must be either on or off which will enable or disable tuples-only mode. If value is omitted the command toggles between regular and tuples-only output. Regular output includes extra information such as column headers, titles, and various footers. In tuples-only mode, only actual table data is shown.
psql

unicode_border_linestyle

Sets the border drawing style for the unicode line style to one of single or double.

unicode_column_linestyle

Sets the column drawing style for the unicode line style to one of single or double.

unicode_header_linestyle

Sets the header drawing style for the unicode line style to one of single or double.

Illustrations of how these different formats look can be seen in the Examples section.

Tip: There are various shortcut commands for \pset. See \a, \C, \f, \H, \t, \T, and \x.

\q or \quit

Quits the psql program. In a script file, only execution of that script is terminated.

\qecho text [ ... ]

This command is identical to \echo except that the output will be written to the query output channel, as set by \o.

\r or \reset

Resets (clears) the query buffer.

\s [ filename ]

Print psql’s command line history to filename. If filename is omitted, the history is written to the standard output (using the pager if appropriate). This command is not available if psql was built without Readline support.

\set [ name [ value [ ... ] ] ]

Sets the psql variable name to value, or if more than one value is given, to the concatenation of all of them. If only one argument is given, the variable is set with an empty value. To unset a variable, use the \unset command.

\set without any arguments displays the names and values of all currently-set psql variables.

Valid variable names can contain letters, digits, and underscores. See the section Variables below for details. Variable names are case-sensitive.

Although you are welcome to set any variable to anything you want, psql treats several variables as special. They are documented in the section about variables.

Note: This command is unrelated to the SQL command SET.

\setenv name [ value ]

Sets the environment variable name to value, or if the value is not supplied, unsets the environment variable. Example:

```
testdb=> \setenv PAGER less
testdb=> \setenv LESS -imx4F
```
\sf[+] function_description

This command fetches and shows the definition of the named function, in the form of a `CREATE OR REPLACE FUNCTION` command. The definition is printed to the current query output channel, as set by `\o`.

The target function can be specified by name alone, or by name and arguments, for example `foo(integer, text)`. The argument types must be given if there is more than one function of the same name.

If `+` is appended to the command name, then the output lines are numbered, with the first line of the function body being line 1.

\sv[+] view_name

This command fetches and shows the definition of the named view, in the form of a `CREATE OR REPLACE VIEW` command. The definition is printed to the current query output channel, as set by `\o`.

If `+` is appended to the command name, then the output lines are numbered from 1.

\t

Toggles the display of output column name headings and row count footer. This command is equivalent to `\pset tuples_only` and is provided for convenience.

\T table_options

Specifies attributes to be placed within the `table` tag in HTML output format. This command is equivalent to `\pset tableattr table_options`.

\timing [ on | off ]

Without parameter, toggles a display of how long each SQL statement takes, in milliseconds. With parameter, sets same.

\unset name

Unsets (deletes) the psql variable `name`.

\w or \write filename
\w or \write \command

Outputs the current query buffer to the file `filename` or pipes it to the shell command `\command`.

\watch [ seconds ]

Repeatedly execute the current query buffer (as `\g` does) until interrupted or the query fails. Wait the specified number of seconds (default 2) between executions. Each query result is displayed with a header that includes the `\pset title` string (if any), the time as of query start, and the delay interval.

\x [ on | off | auto ]

Sets or toggles expanded table formatting mode. As such it is equivalent to `\pset expanded`.

\z [ pattern ]

Lists tables, views and sequences with their associated access privileges. If a `pattern` is specified, only tables, views and sequences whose names match the pattern are listed.

This is an alias for `\dp` (“display privileges”).
\![ \text{command} \]  

Escapes to a separate shell or executes the shell command \text{command}. The arguments are not further interpreted; the shell will see them as-is. In particular, the variable substitution rules and backslash escapes do not apply.

\?[ \text{topic} \]  

Shows help information. The optional \text{topic} parameter (defaulting to \text{commands}) selects which part of psql is explained: \text{commands} describes psql’s backslash commands; \text{options} describes the command-line options that can be passed to psql; and \text{variables} shows help about psql configuration variables.

**Patterns**

The various \text{\d} commands accept a \text{pattern} parameter to specify the object name(s) to be displayed. In the simplest case, a pattern is just the exact name of the object. The characters within a pattern are normally folded to lower case, just as in SQL names; for example, \text{\dt FOO} will display the table named \text{foo}. As in SQL names, placing double quotes around a pattern stops folding to lower case. Should you need to include an actual double quote character in a pattern, write it as a pair of double quotes within a double-quote sequence; again this is in accord with the rules for SQL quoted identifiers. For example, \text{\dt "FOO""BAR"} will display the table named \text{FOO"BAR} (not \text{foo"bar}). Unlike the normal rules for SQL names, you can put double quotes around just part of a pattern, for instance \text{\dt FOO"FOO"BAR} will display the table named \text{fooFOObar}.

Whenever the \text{pattern} parameter is omitted completely, the \text{\d} commands display all objects that are visible in the current schema search path — this is equivalent to using \text{*} as the pattern. (An object is said to be \text{visible} if its containing schema is in the search path and no object of the same kind and name appears earlier in the search path. This is equivalent to the statement that the object can be referenced by name without explicit schema qualification.) To see all objects in the database regardless of visibility, use \text{*.*} as the pattern.

Within a pattern, \text{*} matches any sequence of characters (including no characters) and \text{?} matches any single character. (This notation is comparable to Unix shell file name patterns.) For example, \text{\dt int*} displays tables whose names begin with \text{int}. But within double quotes, \text{*} and \text{?} lose these special meanings and are just matched literally.

A pattern that contains a dot (\text{.}) is interpreted as a schema name pattern followed by an object name pattern. For example, \text{\dt foo.*bar*} displays all tables whose table name includes \text{bar} that are in schemas whose schema name starts with \text{foo}. When no dot appears, then the pattern matches only objects that are visible in the current schema search path. Again, a dot within double quotes loses its special meaning and is matched literally.

Advanced users can use regular-expression notations such as character classes, for example \text{[0-9]} to match any digit. All regular expression special characters work as specified in Section 9.7.3, except for \text{.}, which is taken as a separator as mentioned above, \text{*} which is translated to the regular-expression notation \text{.*}, \text{?} which is translated to \text{.}, and \text{\$} which is matched literally. You can emulate these pattern characters at need by writing \text{?} for \text{.}, \text{(R*)} for \text{R*}, or \text{(R\text{\d})} for \text{R\d}. \text{\$} is not needed as a regular-expression character since the pattern must match the whole name, unlike the usual interpretation of regular expressions (in other words, \text{\$} is automatically appended to your pattern). Write \text{*} at the beginning and/or end if you don’t wish the pattern to be anchored. Note that within double quotes, all regular expression special characters lose their special meanings and are matched literally. Also, the regular expression special characters are matched literally in operator name patterns (i.e., the argument of \text{\do}).
Advanced Features

Variables

psql provides variable substitution features similar to common Unix command shells. Variables are simply name/value pairs, where the value can be any string of any length. The name must consist of letters (including non-Latin letters), digits, and underscores.

To set a variable, use the psql meta-command \set. For example,

```
psql (testdb) SQL> \set foo bar
```

sets the variable `foo` to the value `bar`. To retrieve the content of the variable, precede the name with a colon, for example:

```
psql (testdb) SQL> \echo :foo
```

This works in both regular SQL commands and meta-commands; there is more detail in SQL Interpolation, below.

If you call \set without a second argument, the variable is set, with an empty string as value. To unset (i.e., delete) a variable, use the command \unset. To show the values of all variables, call \set without any argument.

**Note:** The arguments of \set are subject to the same substitution rules as with other commands. Thus you can construct interesting references such as \set :foo 'something' and get “soft links” or “variable variables” of Perl or PHP fame, respectively. Unfortunately (or fortunately?), there is no way to do anything useful with these constructs. On the other hand, \set bar :foo is a perfectly valid way to copy a variable.

A number of these variables are treated specially by psql. They represent certain option settings that can be changed at run time by altering the value of the variable, or in some cases represent changeable state of psql. Although you can use these variables for other purposes, this is not recommended, as the program behavior might grow really strange really quickly. By convention, all specially treated variables’ names consist of all upper-case ASCII letters (and possibly digits and underscores). To ensure maximum compatibility in the future, avoid using such variable names for your own purposes. A list of all specially treated variables follows.

**AUTOcommit**

When on (the default), each SQL command is automatically committed upon successful completion. To postpone commit in this mode, you must enter a `BEGIN` or `START TRANSACTION` SQL command. When off or unset, SQL commands are not committed until you explicitly issue `COMMIT` or `END`. The autocommit-off mode works by issuing an implicit `BEGIN` for you, just before any command that is not already in a transaction block and is not itself a `BEGIN` or other transaction-control command, nor a command that cannot be executed inside a transaction block (such as `VACUUM`).

**Note:** In autocommit-off mode, you must explicitly abandon any failed transaction by entering `ABORT` or `ROLLBACK`. Also keep in mind that if you exit the session without committing, your work will be lost.
**Note:** The autocommit-on mode is PostgreSQL's traditional behavior, but autocommit-off is closer to the SQL spec. If you prefer autocommit-off, you might wish to set it in the system-wide `psqlrc` file or your `~/.psqlrc` file.

**COMP_KEYWORD_CASE**

Determines which letter case to use when completing an SQL key word. If set to `lower` or `upper`, the completed word will be in lower or upper case, respectively. If set to `preserve-lower` or `preserve-upper` (the default), the completed word will be in the case of the word already entered, but words being completed without anything entered will be in lower or upper case, respectively.

**DBNAME**

The name of the database you are currently connected to. This is set every time you connect to a database (including program start-up), but can be unset.

**ECHO**

If set to `all`, all nonempty input lines are printed to standard output as they are read. (This does not apply to lines read interactively.) To select this behavior on program start-up, use the switch `-a`. If set to `queries`, `psql` prints each query to standard output as it is sent to the server. The switch for this is `-e`. If set to `errors`, then only failed queries are displayed on standard error output. The switch for this is `-b`. If unset, or if set to `none` (or any other value than those above) then no queries are displayed.

**ECHO_HIDDEN**

When this variable is set to `on` and a backslash command queries the database, the query is first shown. This feature helps you to study PostgreSQL internals and provide similar functionality in your own programs. (To select this behavior on program start-up, use the switch `-E`.) If you set the variable to the value `noexec`, the queries are just shown but are not actually sent to the server and executed.

**ENCODING**

The current client character set encoding.

**FETCH_COUNT**

If this variable is set to an integer value > 0, the results of `SELECT` queries are fetched and displayed in groups of that many rows, rather than the default behavior of collecting the entire result set before display. Therefore only a limited amount of memory is used, regardless of the size of the result set. Settings of 100 to 1000 are commonly used when enabling this feature. Keep in mind that when using this feature, a query might fail after having already displayed some rows.

**Tip:** Although you can use any output format with this feature, the default `aligned` format tends to look bad because each group of `FETCH_COUNT` rows will be formatted separately, leading to varying column widths across the row groups. The other output formats work better.

**HISTCONTROL**

If this variable is set to `ignorespace`, lines which begin with a space are not entered into the history list. If set to a value of `ignoredups`, lines matching the previous history line are not
entered. A value of ignoreboth combines the two options. If unset, or if set to none (or any other value than those above), all lines read in interactive mode are saved on the history list.

**Note:** This feature was shamelessly plagiarized from Bash.

### HISTFILE

The file name that will be used to store the history list. The default value is `~/.psql_history`. For example, putting:

```
\set HISTFILE ~/.psql_history - :DBNAME
```

in `~/.psqlrc` will cause psql to maintain a separate history for each database.

**Note:** This feature was shamelessly plagiarized from Bash.

### HISTSIZE

The number of commands to store in the command history. The default value is 500.

**Note:** This feature was shamelessly plagiarized from Bash.

### HOST

The database server host you are currently connected to. This is set every time you connect to a database (including program start-up), but can be unset.

### IGNOREEOF

If unset, sending an EOF character (usually `Control+D`) to an interactive session of psql will terminate the application. If set to a numeric value, that many EOF characters are ignored before the application terminates. If the variable is set but has no numeric value, the default is 10.

**Note:** This feature was shamelessly plagiarized from Bash.

### LASTOID

The value of the last affected OID, as returned from an `INSERT` or `\lo_import` command. This variable is only guaranteed to be valid until after the result of the next SQL command has been displayed.

### ON_ERROR_ROLLBACK

When set to `on`, if a statement in a transaction block generates an error, the error is ignored and the transaction continues. When set to `interactive`, such errors are only ignored in interactive sessions, and not when reading script files. When unset or set to `off`, a statement in a transaction block that generates an error aborts the entire transaction. The error rollback mode works by issuing an implicit `SAVEPOINT` for you, just before each command that is in a transaction block, and then rolling back to the savepoint if the command fails.

### ON_ERROR_STOP

By default, command processing continues after an error. When this variable is set to `on`, processing will instead stop immediately. In interactive mode, psql will return to the command prompt;
otherwise, psql will exit, returning error code 3 to distinguish this case from fatal error conditions, which are reported using error code 1. In either case, any currently running scripts (the top-level script, if any, and any other scripts which it may have in invoked) will be terminated immediately. If the top-level command string contained multiple SQL commands, processing will stop with the current command.

PORT
The database server port to which you are currently connected. This is set every time you connect to a database (including program start-up), but can be unset.

PROMPT1
PROMPT2
PROMPT3
These specify what the prompts psql issues should look like. See Prompting below.

QUIET
Setting this variable to on is equivalent to the command line option -q. It is probably not too useful in interactive mode.

SHOW_CONTEXT
This variable can be set to the values never, errors, or always to control whether CONTEXT fields are displayed in messages from the server. The default is errors (meaning that context will be shown in error messages, but not in notice or warning messages). This setting has no effect when VERBOSITY is set to terse. (See also \errverbose, for use when you want a verbose version of the error you just got.)

SINGLELINE
Setting this variable to on is equivalent to the command line option -S.

SINGLESTEP
Setting this variable to on is equivalent to the command line option -s.

USER
The database user you are currently connected as. This is set every time you connect to a database (including program start-up), but can be unset.

VERBOSITY
This variable can be set to the values default, verbose, or terse to control the verbosity of error reports. (See also \errverbose, for use when you want a verbose version of the error you just got.)

SQL Interpolation
A key feature of psql variables is that you can substitute (“interpolate”) them into regular SQL statements, as well as the arguments of meta-commands. Furthermore, psql provides facilities for ensuring that variable values used as SQL literals and identifiers are properly quoted. The syntax for interpolating a value without any quoting is to prepend the variable name with a colon (:). For example,

testdb=> \set foo 'my_table'
testdb=> SELECT * FROM :foo;

would query the table my_table. Note that this may be unsafe: the value of the variable is copied literally, so it can contain unbalanced quotes, or even backslash commands. You must make sure that it makes sense where you put it.
When a value is to be used as an SQL literal or identifier, it is safest to arrange for it to be quoted. To quote the value of a variable as an SQL literal, write a colon followed by the variable name in single quotes. To quote the value as an SQL identifier, write a colon followed by the variable name in double quotes. These constructs deal correctly with quotes and other special characters embedded within the variable value. The previous example would be more safely written this way:

```
testdb=> \set foo 'my_table'
```
```
testdb=> SELECT * FROM :"foo";
```

Variable interpolation will not be performed within quoted SQL literals and identifiers. Therefore, a construction such as `':foo'` doesn’t work to produce a quoted literal from a variable’s value (and it would be unsafe if it did work, since it wouldn’t correctly handle quotes embedded in the value).

One example use of this mechanism is to copy the contents of a file into a table column. First load the file into a variable and then interpolate the variable’s value as a quoted string:

```
testdb=> \set content 'cat my_file.txt'
```
```
testdb=> INSERT INTO my_table VALUES (:'content');
```

(Note that this still won’t work if `my_file.txt` contains NUL bytes. psql does not support embedded NUL bytes in variable values.)

Since colons can legally appear in SQL commands, an apparent attempt at interpolation (that is, `:name`, `:'name'`, or `:"name"`) is not replaced unless the named variable is currently set. In any case, you can escape a colon with a backslash to protect it from substitution.

The colon syntax for variables is standard SQL for embedded query languages, such as ECPG. The colon syntaxes for array slices and type casts are PostgreSQL extensions, which can sometimes conflict with the standard usage. The colon-quote syntax for escaping a variable’s value as an SQL literal or identifier is a psql extension.

**Prompting**

The prompts psql issues can be customized to your preference. The three variables `PROMPT1`, `PROMPT2`, and `PROMPT3` contain strings and special escape sequences that describe the appearance of the prompt. Prompt 1 is the normal prompt that is issued when psql requests a new command. Prompt 2 is issued when more input is expected during command entry, for example because the command was not terminated with a semicolon or a quote was not closed. Prompt 3 is issued when you are running an SQL `COPY FROM STDIN` command and you need to type in a row value on the terminal.

The value of the selected prompt variable is printed literally, except where a percent sign (%) is encountered. Depending on the next character, certain other text is substituted instead. Defined substitutions are:

```
%M
```

The full host name (with domain name) of the database server, or `[local]` if the connection is over a Unix domain socket, or `[local:/dir/name]`, if the Unix domain socket is not at the compiled in default location.

```
%m
```

The host name of the database server, truncated at the first dot, or `[local]` if the connection is over a Unix domain socket.
The port number at which the database server is listening.

The database session user name. (The expansion of this value might change during a database session as the result of the command `SET SESSION AUTHORIZATION`.)

The name of the current database.

Like `%`, but the output is ~ (tilde) if the database is your default database.

If the session user is a database superuser, then a #, otherwise a >. (The expansion of this value might change during a database session as the result of the command `SET SESSION AUTHORIZATION`.)

The process ID of the backend currently connected to.

In prompt 1 normally =, but ^ if in single-line mode, or ! if the session is disconnected from the database (which can happen if `\connect` fails). In prompt 2 `%R` is replaced by a character that depends on why `psql` expects more input: = if the command simply wasn’t terminated yet, but * if there is an unfinished /* ... */ comment, a single quote if there is an unfinished quoted string, a double quote if there is an unfinished quoted identifier, a dollar sign if there is an unfinished dollar-quoted string, or ( if there is an unmatched left parenthesis. In prompt 3 `%R` doesn’t produce anything.

Transaction status: an empty string when not in a transaction block, or * when in a transaction block, or ! when in a failed transaction block, or ? when the transaction state is indeterminate (for example, because there is no connection).

The line number inside the current statement, starting from 1.

The character with the indicated octal code is substituted.

The value of the `psql` variable `name`. See the section Variables for details.

The output of `command`, similar to ordinary “back-tick” substitution.

Prompts can contain terminal control characters which, for example, change the color, background, or style of the prompt text, or change the title of the terminal window. In order for the line editing features of Readline to work properly, these non-printing control characters must be designated as invisible by surrounding them with `%[ and `%]. Multiple pairs of these can occur within the prompt. For example:

testdb=> \set PROMPT1 ‘%[033[1;33;40m]%n@%/%R%[033[0m]%# ’
results in a boldfaced (1;) yellow-on-black (33;40) prompt on VT100-compatible, color-capable terminals.

To insert a percent sign into your prompt, write %% . The default prompts are ‘%/%R%# ’ for prompts 1 and 2, and ‘>> ’ for prompt 3.

Note: This feature was shamelessly plagiarized from tcsh.

Command-Line Editing

psql supports the Readline library for convenient line editing and retrieval. The command history is automatically saved when psql exits and is reloaded when psql starts up. Tab-completion is also supported, although the completion logic makes no claim to be an SQL parser. The queries generated by tab-completion can also interfere with other SQL commands, e.g. SET TRANSACTION ISOLATION LEVEL. If for some reason you do not like the tab completion, you can turn it off by putting this in a file named .inputrc in your home directory:

```
$if psql
set disable-completion on
$endif
```

(This is not a psql but a Readline feature. Read its documentation for further details.)

Environment

COLUMNS

If \set columns is zero, controls the width for the wrapped format and width for determining if wide output requires the pager or should be switched to the vertical format in expanded auto mode.

PAGER

If the query results do not fit on the screen, they are piped through this command. Typical values are more or less. The default is platform-dependent. Use of the pager can be disabled by setting PAGER to empty, or by using pager-related options of the \psset command.

PGDATABASE
PGHOST
PGPORT
PGUSER

Default connection parameters (see Section 32.14).

PSQL_EDITOR
EDITOR
VISUAL

Editor used by the \e, \ef, and \ev commands. These variables are examined in the order listed; the first that is set is used.

The built-in default editors are vi on Unix systems and notepad.exe on Windows systems.
When \e, \ef, or \ev is used with a line number argument, this variable specifies the command-line argument used to pass the starting line number to the user’s editor. For editors such as Emacs or vi, this is a plus sign. Include a trailing space in the value of the variable if there needs to be space between the option name and the line number. Examples:

```
PSQL_EDITOR_LINENUMBER_ARG='+'
PSQL_EDITOR_LINENUMBER_ARG='--line '  
```

The default is + on Unix systems (corresponding to the default editor vi, and useful for many other common editors); but there is no default on Windows systems.

**PSQL_HISTORY**

Alternative location for the command history file. Tilde (~) expansion is performed.

**PSQLRC**

Alternative location of the user’s .psqlrc file. Tilde (~) expansion is performed.

**SHELL**

Command executed by the \! command.

**TMPDIR**

Directory for storing temporary files. The default is /tmp.

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).

## Files

psqlrc and ~/.psqlrc

Unless it is passed an -X option, psql attempts to read and execute commands from the system-wide startup file (psqlrc) and then the user’s personal startup file (~/.psqlrc), after connecting to the database but before accepting normal commands. These files can be used to set up the client and/or the server to taste, typically with \set and SET commands.

The system-wide startup file is named psqlrc and is sought in the installation’s “system configuration” directory, which is most reliably identified by running pg_config --sysconfdir. By default this directory will be ../etc/ relative to the directory containing the PostgreSQL executables. The name of this directory can be set explicitly via the PGSYSCONFDIR environment variable.

The user’s personal startup file is named .psqlrc and is sought in the invoking user’s home directory. On Windows, which lacks such a concept, the personal startup file is named %APPDATA%\postgresql\psqlrc.conf. The location of the user’s startup file can be set explicitly via the PSQLRC environment variable.

Both the system-wide startup file and the user’s personal startup file can be made psql-version-specific by appending a dash and the PostgreSQL major or minor release number to the file name, for example ~/.psqlrc-9.2 or ~/.psqlrc-9.2.5. The most specific version-matching file will be read in preference to a non-version-specific file.

**.psql_history**

The command-line history is stored in the file ~/.psql_history, or %APPDATA%\postgresql\psql_history on Windows.
The location of the history file can be set explicitly via the \$PSQL_HISTORY environment variable.

Notes

- psql works best with servers of the same or an older major version. Backslash commands are particularly likely to fail if the server is of a newer version than psql itself. However, backslash commands of the \family should work with servers of versions back to 7.4, though not necessarily with servers newer than psql itself. The general functionality of running SQL commands and displaying query results should also work with servers of a newer major version, but this cannot be guaranteed in all cases.

If you want to use psql to connect to several servers of different major versions, it is recommended that you use the newest version of psql. Alternatively, you can keep around a copy of psql from each major version and be sure to use the version that matches the respective server. But in practice, this additional complication should not be necessary.

- Before PostgreSQL 9.6, the \texttt{-c} option implied \texttt{-x(\text--no-psqlrc)}; this is no longer the case.

- Before PostgreSQL 8.4, psql allowed the first argument of a single-letter backslash command to start directly after the command, without intervening whitespace. Now, some whitespace is required.

Notes for Windows Users

psql is built as a “console application”. Since the Windows console windows use a different encoding than the rest of the system, you must take special care when using 8-bit characters within psql. If psql detects a problematic console code page, it will warn you at startup. To change the console code page, two things are necessary:

- Set the code page by entering \texttt{cmd.exe /c chcp 1252}. (1252 is a code page that is appropriate for German; replace it with your value.) If you are using Cygwin, you can put this command in /etc/profile.

- Set the console font to \texttt{Lucida Console}, because the raster font does not work with the ANSI code page.

Examples

The first example shows how to spread a command over several lines of input. Notice the changing prompt:

```
testdb=> CREATE TABLE my_table (  
  testdb(>  first integer not null default 0,  
  testdb(>  second text)  
  testdb=> ;  
CREATE TABLE
```

Now look at the table definition again:
```sql
testdb=> \d my_table
Table "my_table"

<table>
<thead>
<tr>
<th>Attribute</th>
<th>Type</th>
<th>Modifier</th>
</tr>
</thead>
<tbody>
<tr>
<td>first</td>
<td>integer</td>
<td>not null default 0</td>
</tr>
<tr>
<td>second</td>
<td>text</td>
<td></td>
</tr>
</tbody>
</table>

Now we change the prompt to something more interesting:

testdb=> \set PROMPT1 '%n@%m %~%R%# '
peter@localhost testdb=>

Let’s assume you have filled the table with data and want to take a look at it:

peter@localhost testdb=> SELECT * FROM my_table;

<table>
<thead>
<tr>
<th>first</th>
<th>second</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>one</td>
</tr>
<tr>
<td>2</td>
<td>two</td>
</tr>
<tr>
<td>3</td>
<td>three</td>
</tr>
<tr>
<td>4</td>
<td>four</td>
</tr>
</tbody>
</table>

(4 rows)

You can display tables in different ways by using the \pset command:

peter@localhost testdb=> \pset border 2
Border style is 2.
peter@localhost testdb=> SELECT * FROM my_table;

<table>
<thead>
<tr>
<th>first</th>
<th>second</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>one</td>
</tr>
<tr>
<td>2</td>
<td>two</td>
</tr>
<tr>
<td>3</td>
<td>three</td>
</tr>
<tr>
<td>4</td>
<td>four</td>
</tr>
</tbody>
</table>

(4 rows)

peter@localhost testdb=> \pset border 0
Border style is 0.
peter@localhost testdb=> SELECT * FROM my_table;

<table>
<thead>
<tr>
<th>first</th>
<th>second</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>one</td>
</tr>
<tr>
<td>2</td>
<td>two</td>
</tr>
<tr>
<td>3</td>
<td>three</td>
</tr>
<tr>
<td>4</td>
<td>four</td>
</tr>
</tbody>
</table>

(4 rows)

peter@localhost testdb=> \pset border 1
Border style is 1.
peter@localhost testdb=> \pset format unaligned
Output format is unaligned.
peter@localhost testdb=> \pset fieldsep ",,"
Field separator is ",,.
peter@localhost testdb=> \pset tuples_only
Showing only tuples.
peter@localhost testdb=> SELECT second, first FROM my_table;
```

1886
one, 1
two, 2
three, 3
four, 4

Alternatively, use the short commands:

```
peter@localhost testdb=> \\a \t \\x
Output format is aligned.
Tuples only is off.
Expanded display is on.
peter@localhost testdb=> SELECT * FROM my_table;
- [ RECORD 1 ] -
first  | 1
second | one
- [ RECORD 2 ] -
first  | 2
second | two
- [ RECORD 3 ] -
first  | 3
second | three
- [ RECORD 4 ] -
first  | 4
second | four
```

When suitable, query results can be shown in a crosstab representation with the `\crosstabview` command:

```
testdb=> SELECT first, second, first > 2 AS gt2 FROM my_table;
first | second | gt2
-------+--------+-----
1     | one    | f
2     | two    | f
3     | three  | t
4     | four   | t
(4 rows)
testdb=> \crosstabview first second
first | one  | two  | three | four
-------+-----+-----+-------+-----
1     | f   |     |       |     
2     |     | f   |       |     
3     |     |     | t     |     
4     |     |     |       | t
(4 rows)
```

This second example shows a multiplication table with rows sorted in reverse numerical order and columns with an independent, ascending numerical order.

```
testdb=> SELECT t1.first as "A", t2.first+100 AS "B", t1.first*(t2.first+100) as "AxB", row_number() over(order by t2.first) AS ord FROM my_table t1 CROSS JOIN my_table t2 ORDER BY 1 DESC
```

```
A | B          | AxB         | ord
---+------------+-------------+-----
101 | 102 | 103 | 104
4 | 404 | 408 | 412 | 416
```

1887
<p>| | | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>303</td>
<td>306</td>
<td>309</td>
<td>312</td>
</tr>
<tr>
<td>2</td>
<td>202</td>
<td>204</td>
<td>206</td>
<td>208</td>
</tr>
<tr>
<td>1</td>
<td>101</td>
<td>102</td>
<td>103</td>
<td>104</td>
</tr>
</tbody>
</table>

(4 rows)
Name

reindexdb — reindex a PostgreSQL database

Synopsis

reindexdb [connection-option...] [option...] [ --schema|--S schema ]... [ --table|--t table ]... [ --index|--i index ]... [ dbname ]

reindexdb [connection-option...] [option...] --all|--a

reindexdb [connection-option...] [option...] --system|--s [ dbname ]

Description

reindexdb is a utility for rebuilding indexes in a PostgreSQL database.

reindexdb is a wrapper around the SQL command REINDEX. There is no effective difference between reindexing databases via this utility and via other methods for accessing the server.

Options

reindexdb accepts the following command-line arguments:

-`a`
-`--all`
  Reindex all databases.

[-d] dbname
[<dbname>=] dbname

  Specifies the name of the database to be reindexed. If this is not specified and `-a` (or `--all`) is not used, the database name is read from the environment variable PGDATABASE. If that is not set, the user name specified for the connection is used.

-`e`
-`--echo`
  Echo the commands that reindexdb generates and sends to the server.

-`i` index
-`--index=` index
  Recreate index only. Multiple indexes can be recreated by writing multiple `-i` switches.

-`q`
-`--quiet`
  Do not display progress messages.
-s  
--system

Reindex database’s system catalogs.

-S schema
--schema=schema

Reindex schema only. Multiple schemas can be reindexed by writing multiple -S switches.

-t table
--table=table

Reindex table only. Multiple tables can be reindexed by writing multiple -t switches.

-v
--verbose

Print detailed information during processing.

-V
--version

Print the reindexdb version and exit.

-?
--help

Show help about reindexdb command line arguments, and exit.

reindexdb also accepts the following command-line arguments for connection parameters:

-h host
--host=host

Specifies the host name of the machine on which the server is running. If the value begins with a
slash, it is used as the directory for the Unix domain socket.

-p port
--port=port

Specifies the TCP port or local Unix domain socket file extension on which the server is listening
for connections.

-U username
--username=username

User name to connect as.

-W
--no-password

Never issue a password prompt. If the server requires password authentication and a password
is not available by other means such as a .pgpass file, the connection attempt will fail. This
option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
--password

Force reindexdb to prompt for a password before connecting to a database.

This option is never essential, since reindexdb will automatically prompt for a password if the
server demands password authentication. However, reindexdb will waste a connection attempt
reindexdb

finding out that the server wants a password. In some cases it is worth typing \(-W\) to avoid the extra connection attempt.

\[---maintenance-db=dbname\]

Specifies the name of the database to connect to discover what other databases should be reindexed. If not specified, the postgres database will be used, and if that does not exist, template1 will be used.

**Environment**

PGDATABASE

PGHOST

PGPORT

PGUSER

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).

**Diagnostics**

In case of difficulty, see REINDEX and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

**Notes**

reindexdb might need to connect several times to the PostgreSQL server, asking for a password each time. It is convenient to have a \(~/.pgpass\) file in such cases. See Section 32.15 for more information.

**Examples**

To reindex the database test:

```
$ reindexdb test
```

To reindex the table foo and the index bar in a database named abcd:

```
$ reindexdb --table foo --index bar abcd
```

**See Also**

REINDEX
vacuumdb

Name

vacuumdb — garbage-collect and analyze a PostgreSQL database

Synopsis

vacuumdb [connection-option...] [option...] [ --table | -t table [( column[,...] )] ] ... [dbname]

vacuumdb [connection-option...] [option...] --all | -a

Description

vacuumdb is a utility for cleaning a PostgreSQL database. vacuumdb will also generate internal statistics used by the PostgreSQL query optimizer.

vacuumdb is a wrapper around the SQL command VACUUM. There is no effective difference between vacuuming and analyzing databases via this utility and via other methods for accessing the server.

Options

vacuumdb accepts the following command-line arguments:

- 
-a
--all
Vacuum all databases.
[-d] dbname
[--dbname=] dbname

Specifies the name of the database to be cleaned or analyzed. If this is not specified and -a (or --all) is not used, the database name is read from the environment variable PGDATABASE. If that is not set, the user name specified for the connection is used.

-e
--echo

Echo the commands that vacuumdb generates and sends to the server.

-f
--full

Perform “full” vacuuming.

-F
--freeze

Aggressively “freeze” tuples.
vacuumdb

-\j njobs
--jobs=njobs

Execute the vacuum or analyze commands in parallel by running \textit{njobs} commands simultaneously. This option reduces the time of the processing but it also increases the load on the database server.

vacuumdb will open \textit{njobs} connections to the database, so make sure your \texttt{max_connections} setting is high enough to accommodate all connections.

Note that using this mode together with the -f (FULL) option might cause deadlock failures if certain system catalogs are processed in parallel.

-q
--quiet

Do not display progress messages.

-t table [ (column [,...]) ]
--table=table [ (column [,...]) ]

Clean or analyze \textit{table} only. Column names can be specified only in conjunction with the --analyze or --analyze-only options. Multiple tables can be vacuumed by writing multiple -t switches.

\textbf{Tip:} If you specify columns, you probably have to escape the parentheses from the shell. (See examples below.)

-v
--verbose

Print detailed information during processing.

-V
--version

Print the vacuumdb version and exit.

-z
--analyze

Also calculate statistics for use by the optimizer.

-Z
--analyze-only

Only calculate statistics for use by the optimizer (no vacuum).

--analyze-in-stages

Only calculate statistics for use by the optimizer (no vacuum), like --analyze-only. Run several (currently three) stages of analyze with different configuration settings, to produce usable statistics faster.

This option is useful to analyze a database that was newly populated from a restored dump or by \texttt{pg\_upgrade}. This option will try to create some statistics as fast as possible, to make the database usable, and then produce full statistics in the subsequent stages.

-?
--help

Show help about vacuumdb command line arguments, and exit.
vacuumdb also accepts the following command-line arguments for connection parameters:

- **-h host**
  - **--host=host**
    - Specifies the host name of the machine on which the server is running. If the value begins with a slash, it is used as the directory for the Unix domain socket.

- **-p port**
  - **--port=port**
    - Specifies the TCP port or local Unix domain socket file extension on which the server is listening for connections.

- **-U username**
  - **--username=username**
    - User name to connect as.

- **-w**
  - **--no-password**
    - Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a `pgpass` file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

- **-W**
  - **--password**
    - Force vacuumdb to prompt for a password before connecting to a database.
    - This option is never essential, since vacuumdb will automatically prompt for a password if the server demands password authentication. However, vacuumdb will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing `-W` to avoid the extra connection attempt.

- **--maintenance-db=dbname**
  - Specifies the name of the database to connect to discover what other databases should be vacuumed. If not specified, the `postgres` database will be used, and if that does not exist, `template1` will be used.

### Environment

- **PGDATABASE**
- **PGHOST**
- **PGPORT**
- **PGUSER**

Default connection parameters

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).
Diagnostics

In case of difficulty, see VACUUM and psql for discussions of potential problems and error messages. The database server must be running at the targeted host. Also, any default connection settings and environment variables used by the libpq front-end library will apply.

Notes

vacuumdb might need to connect several times to the PostgreSQL server, asking for a password each time. It is convenient to have a ~/.pgpass file in such cases. See Section 32.15 for more information.

Examples

To clean the database test:

$ vacuumdb test

To clean and analyze for the optimizer a database named bigdb:

$ vacuumdb --analyze bigdb

To clean a single table foo in a database named xyzzy, and analyze a single column bar of the table for the optimizer:

$ vacuumdb --analyze --verbose --table 'foo(bar)' xyzzy

See Also

VACUUM
III. PostgreSQL Server Applications

This part contains reference information for PostgreSQL server applications and support utilities. These commands can only be run usefully on the host where the database server resides. Other utility programs are listed in Reference II, *PostgreSQL Client Applications*.
initdb

Name

initdb — create a new PostgreSQL database cluster

Synopsis

initdb [option...] [--pgdata=-] directory

Description

initdb creates a new PostgreSQL database cluster. A database cluster is a collection of databases that are managed by a single server instance.

Creating a database cluster consists of creating the directories in which the database data will live, generating the shared catalog tables (tables that belong to the whole cluster rather than to any particular database), and creating the template1 and postgres databases. When you later create a new database, everything in the template1 database is copied. (Therefore, anything installed in template1 is automatically copied into each database created later.) The postgres database is a default database meant for use by users, utilities and third party applications.

Although initdb will attempt to create the specified data directory, it might not have permission if the parent directory of the desired data directory is root-owned. To initialize in such a setup, create an empty data directory as root, then use chown to assign ownership of that directory to the database user account, then su to become the database user to run initdb.

initdb must be run as the user that will own the server process, because the server needs to have access to the files and directories that initdb creates. Since the server cannot be run as root, you must not run initdb as root either. (It will in fact refuse to do so.)

initdb initializes the database cluster’s default locale and character set encoding. The character set encoding, collation order (LC_COLLATE) and character set classes (LC_CTYPE, e.g. upper, lower, digit) can be set separately for a database when it is created. initdb determines those settings for the template1 database, which will serve as the default for all other databases.

To alter the default collation order or character set classes, use the --lc-collate and --lc-ctype options. Collation orders other than C or POSIX also have a performance penalty. For these reasons it is important to choose the right locale when running initdb.

The remaining locale categories can be changed later when the server is started. You can also use --locale to set the default for all locale categories, including collation order and character set classes. All server locale values (lc_*) can be displayed via SHOW ALL. More details can be found in Section 23.1.

To alter the default encoding, use the --encoding. More details can be found in Section 23.3.
Options

-A authmethod
--auth=authmethod

This option specifies the authentication method for local users used in pg_hba.conf (host and local lines). Do not use trust unless you trust all local users on your system. trust is the default for ease of installation.

--auth-host=authmethod

This option specifies the authentication method for local users via TCP/IP connections used in pg_hba.conf (host lines).

--auth-local=authmethod

This option specifies the authentication method for local users via Unix-domain socket connections used in pg_hba.conf (local lines).

-D directory
--pgdata=directory

This option specifies the directory where the database cluster should be stored. This is the only information required by initdb, but you can avoid writing it by setting the PGDATA environment variable, which can be convenient since the database server (postgres) can find the database directory later by the same variable.

-E encoding
--encoding=encoding

Selects the encoding of the template database. This will also be the default encoding of any database you create later, unless you override it there. The default is derived from the locale, or SQL_ASCII if that does not work. The character sets supported by the PostgreSQL server are described in Section 23.3.1.

-k
--data-checksums

Use checksums on data pages to help detect corruption by the I/O system that would otherwise be silent. Enabling checksums may incur a noticeable performance penalty. This option can only be set during initialization, and cannot be changed later. If set, checksums are calculated for all objects, in all databases.

--locale=locale

Sets the default locale for the database cluster. If this option is not specified, the locale is inherited from the environment that initdb runs in. Locale support is described in Section 23.1.

--lc-collate=locale
--lc-type=locale
--lc-messages=locale
--lc-monetary=locale
--lc-numeric=locale
--lc-time=locale

Like --locale, but only sets the locale in the specified category.

--no-locale

Equivalent to --locale=C.
By default, `initdb` will wait for all files to be written safely to disk. This option causes `initdb` to return without waiting, which is faster, but means that a subsequent operating system crash can leave the data directory corrupt. Generally, this option is useful for testing, but should not be used when creating a production installation.

`--pwfile=filename`

Makes `initdb` read the database superuser’s password from a file. The first line of the file is taken as the password.

`--sync-only`

Safely write all database files to disk and exit. This does not perform any of the normal `initdb` operations.

`--text-search-config=CFG`

Sets the default text search configuration. See `default_text_search_config` for further information.

`--username=username`

Selects the user name of the database superuser. This defaults to the name of the effective user running `initdb`. It is really not important what the superuser’s name is, but one might choose to keep the customary name postgres, even if the operating system user’s name is different.

`--xlogdir=directory`

This option specifies the directory where the transaction log should be stored.

Other, less commonly used, options are also available:

`--debug`

Print debugging output from the bootstrap backend and a few other messages of lesser interest for the general public. The bootstrap backend is the program `initdb` uses to create the catalog tables. This option generates a tremendous amount of extremely boring output.

`--xlogdir=directory`

Specifies where `initdb` should find its input files to initialize the database cluster. This is normally not necessary. You will be told if you need to specify their location explicitly.
By default, when `initdb` determines that an error prevented it from completely creating the database cluster, it removes any files it might have created before discovering that it cannot finish the job. This option inhibits tidying-up and is thus useful for debugging.

Other options:

- `-v`
- `--version`
  
  Print the initdb version and exit.

- `--help`

  Show help about initdb command line arguments, and exit.

**Environment**

**PGDATA**

Specifies the directory where the database cluster is to be stored; can be overridden using the `-D` option.

**TZ**

Specifies the default time zone of the created database cluster. The value should be a full time zone name (see Section 8.5.3).

This utility, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14).

**Notes**

`initdb` can also be invoked via `pg_ctl initdb`.

**See Also**

`pg_ctl`, `postgres`
**pg_archivecleanup**

**Name**

`pg_archivecleanup` — clean up PostgreSQL WAL archive files

**Synopsis**

`pg_archivecleanup [option...] archivelocation oldestkeptwalfile`

**Description**

`pg_archivecleanup` is designed to be used as an `archive_cleanup_command` to clean up WAL file archives when running as a standby server (see Section 26.2). `pg_archivecleanup` can also be used as a standalone program to clean WAL file archives.

To configure a standby server to use `pg_archivecleanup`, put this into its `recovery.conf` configuration file:

```
archive_cleanup_command = 'pg_archivecleanup archivelocation %r'
```

where `archivelocation` is the directory from which WAL segment files should be removed.

When used within `archive_cleanup_command`, all WAL files logically preceding the value of the `%r` argument will be removed from `archivelocation`. This minimizes the number of files that need to be retained, while preserving crash-restart capability. Use of this parameter is appropriate if the `archivelocation` is a transient staging area for this particular standby server, but not when the `archivelocation` is intended as a long-term WAL archive area, or when multiple standby servers are recovering from the same archive location.

When used as a standalone program all WAL files logically preceding the `oldestkeptwalfile` will be removed from `archivelocation`. In this mode, if you specify a `.partial` or `.backup` file name, then only the file prefix will be used as the `oldestkeptwalfile`. This treatment of `.backup` file name allows you to remove all WAL files archived prior to a specific base backup without error. For example, the following example will remove all files older than WAL file name `000000010000003700000010`:

```
pg_archivecleanup -d archive 00000001000000370000000010.000000020.backup
```

`pg_archivecleanup` assumes that `archivelocation` is a directory readable and writable by the server-owning user.
Options

pg_archivecleanup accepts the following command-line arguments:

- **d**
  Print lots of debug logging output on stderr.

- **n**
  Print the names of the files that would have been removed on stdout (performs a dry run).

- **V**
  --version
  Print the pg_archivecleanup version and exit.

- **x extension**
  Provide an extension that will be stripped from all file names before deciding if they should be deleted. This is typically useful for cleaning up archives that have been compressed during storage, and therefore have had an extension added by the compression program. For example:
  - **x .gz**.

- **?**
  --help
  Show help about pg_archivecleanup command line arguments, and exit.

Notes

pg_archivecleanup is designed to work with PostgreSQL 8.0 and later when used as a standalone utility, or with PostgreSQL 9.0 and later when used as an archive cleanup command.

pg_archivecleanup is written in C and has an easy-to-modify source code, with specifically designated sections to modify for your own needs

Examples

On Linux or Unix systems, you might use:

```
archive_cleanup_command = 'pg_archivecleanup -d /mnt/standby/archive %r 2>>cleanup.log'
```

where the archive directory is physically located on the standby server, so that the archive command is accessing it across NFS, but the files are local to the standby. This will:

- produce debugging output in cleanup.log
- remove no-longer-needed files from the archive directory

See Also

pg_standby
**pg_controldata**

**Name**

`pg_controldata` — display control information of a PostgreSQL database cluster

**Synopsis**

```
pg_controldata [option] [[-D] datadir]
```

**Description**

`pg_controldata` prints information initialized during `initdb`, such as the catalog version. It also shows information about write-ahead logging and checkpoint processing. This information is cluster-wide, and not specific to any one database.

This utility can only be run by the user who initialized the cluster because it requires read access to the data directory. You can specify the data directory on the command line, or use the environment variable `PGDATA`. This utility supports the options `-V` and `--version`, which print the `pg_controldata` version and exit. It also supports options `--?` and `--help`, which output the supported arguments.

**Environment**

`PGDATA`

Default data directory location
**pg_ctl**

**Name**

pg_ctl — initialize, start, stop, or control a PostgreSQL server

**Synopsis**

`pg_ctl init[db] [-s][-D datadir][-o initdb-options]`

`pg_ctl start [-w][-t seconds][-s][-D datadir][-l filename][-o options][-p path][-c]`

`pg_ctl stop [-W][-t seconds][-s][-D datadir][-m s[mart]|f[ast]|i[mediate]]`

`pg_ctl restart [-w][-t seconds][-s][-D datadir][-c][-m s[mart]|f[ast]|i[mediate]][-o options]`

`pg_ctl reload [-s][-D datadir]`

`pg_ctl status [-D datadir]`

`pg_ctl promote [-s][-D datadir]`

`pg_ctl kill signal_name process_id`

`pg_ctl register [-N servicename][-U username][-P password][-D datadir][-S a[uto]|d[emand]][-w][-t seconds][-s][-o options]`

`pg_ctl unregister [-N servicename]`

**Description**

pg_ctl is a utility for initializing a PostgreSQL database cluster, starting, stopping, or restarting the PostgreSQL database server (postgresql), or displaying the status of a running server. Although the server can be started manually, pg_ctl encapsulates tasks such as redirecting log output and properly detaching from the terminal and process group. It also provides convenient options for controlled shutdown.

The `init` or `initdb` mode creates a new PostgreSQL database cluster. A database cluster is a collection of databases that are managed by a single server instance. This mode invokes the `initdb` command. See `initdb` for details.

In `start` mode, a new server is launched. The server is started in the background, and its standard input is attached to `/dev/null` (or `nul` on Windows). On Unix-like systems, by default, the server’s
standard output and standard error are sent to pg_ctl’s standard output (not standard error). The standard output of pg_ctl should then be redirected to a file or piped to another process such as a log rotating program like rotatelogs; otherwise postgres will write its output to the controlling terminal (from the background) and will not leave the shell’s process group. On Windows, by default the server’s standard output and standard error are sent to the terminal. These default behaviors can be changed by using \(-l\) to append the server’s output to a log file. Use of either \(-l\) or output redirection is recommended.

In \texttt{stop} mode, the server that is running in the specified data directory is shut down. Three different shutdown methods can be selected with the \texttt{-m} option. “Smart” mode waits for all active clients to disconnect and any online backup to finish. If the server is in hot standby, recovery and streaming replication will be terminated once all clients have disconnected. “Fast” mode (the default) does not wait for clients to disconnect and will terminate an online backup in progress. All active transactions are rolled back and clients are forcibly disconnected, then the server is shut down. “Immediate” mode will abort all server processes immediately, without a clean shutdown. This will lead to a crash-recovery run on the next restart.

\texttt{restart} mode effectively executes a stop followed by a start. This allows changing the \texttt{postgres} command-line options. \texttt{restart} might fail if relative paths specified were specified on the command-line during server start.

\texttt{reload} mode simply sends the \texttt{postgres} process a SIGHUP signal, causing it to reread its configuration files (\texttt{postgresql.conf}, \texttt{pg_hba.conf}, etc.). This allows changing of configuration-file options that do not require a complete restart to take effect.

\texttt{status} mode checks whether a server is running in the specified data directory. If it is, the PID and the command line options that were used to invoke it are displayed. If the server is not running, the process returns an exit status of 3. If an accessible data directory is not specified, the process returns an exit status of 4.

In \texttt{promote} mode, the standby server that is running in the specified data directory is commanded to exit recovery and begin read-write operations.

\texttt{kill} mode allows you to send a signal to a specified process. This is particularly valuable for Microsoft Windows which does not have a kill command. Use \texttt{--help} to see a list of supported signal names.

\texttt{register} mode allows you to register a system service on Microsoft Windows. The \texttt{-S} option allows selection of service start type, either “auto” (start service automatically on system startup) or “demand” (start service on demand).

\texttt{unregister} mode allows you to unregister a system service on Microsoft Windows. This undoes the effects of the \texttt{register} command.

\textbf{Options}

\texttt{-c}

\texttt{--core-file}

Attempt to allow server crashes to produce core files, on platforms where this is possible, by lifting any soft resource limit placed on core files. This is useful in debugging or diagnosing problems by allowing a stack trace to be obtained from a failed server process.
-D datadir
--pgdata datadir

Specifies the file system location of the database configuration files. If this is omitted, the environment variable PGDATA is used.

-l filename
--log filename

Append the server log output to filename. If the file does not exist, it is created. The umask is set to 077, so access to the log file is disallowed to other users by default.

-m mode
--mode mode

Specifies the shutdown mode. mode can be smart, fast, or immediate, or the first letter of one of these three. If this is omitted, fast is used.

-o options

Specifies options to be passed directly to the postgres command; multiple option invocations are appended.

The options should usually be surrounded by single or double quotes to ensure that they are passed through as a group.

-o initdb-options

Specifies options to be passed directly to the initdb command.

The options should usually be surrounded by single or double quotes to ensure that they are passed through as a group.

-p path

Specifies the location of the postgres executable. By default the postgres executable is taken from the same directory as pg_ctl, or failing that, the hard-wired installation directory. It is not necessary to use this option unless you are doing something unusual and get errors that the postgres executable was not found.

In init mode, this option analogously specifies the location of the initdb executable.

-s
--silent

Print only errors, no informational messages.

-t
--timeout

The maximum number of seconds to wait when waiting for startup or shutdown to complete. Defaults to the value of the PGCTLTIMEOUT environment variable or, if not set, to 60 seconds.

-V
--version

Print the pg_ctl version and exit.

-w

Wait for the startup or shutdown to complete. Waiting is the default option for shutdowns, but not startups. When waiting for startup, pg_ctl repeatedly attempts to connect to the server. When waiting for shutdown, pg_ctl waits for the server to remove its PID file. This option allows the entry of an SSL passphrase on startup. pg_ctl returns an exit code based on the success of the startup or shutdown.
pg_ctl

-W
Do not wait for startup or shutdown to complete. This is the default for start and restart modes.

-?
--help
Show help about pg_ctl command line arguments, and exit.

Options for Windows

-e source
Name of the event source for pg_ctl to use for logging to the event log when running as a Windows service. The default is PostgreSQL. Note that this only controls the logging from pg_ctl itself; once started, the server will use the event source specified by event_source. Should the server fail during early startup, it might also log using the default event source PostgreSQL.

-N servicename
Name of the system service to register. The name will be used as both the service name and the display name.

-P password
Password for the user to start the service.

-S start-type
Start type of the system service to register. start-type can be auto, or demand, or the first letter of one of these two. If this is omitted, auto is used.

-U username
User name for the user to start the service. For domain users, use the format DOMAIN\username.

Environment

PGCTLTIMEOUT
Default limit on the number of seconds to wait when waiting for startup or shutdown to complete. If not set, the default is 60 seconds.

PGDATA
Default data directory location.

pg_ctl, like most other PostgreSQL utilities, also uses the environment variables supported by libpq (see Section 32.14). For additional server variables, see postgres.

Files

postmaster.pid
The existence of this file in the data directory is used to help pg_ctl determine if the server is currently running.
If this file exists in the data directory, `pg_ctl` (in \texttt{restart} mode) will pass the contents of the file as options to postgres, unless overridden by the \texttt{-o} option. The contents of this file are also displayed in \texttt{status} mode.

**Examples**

**Starting the Server**

To start the server:

\$ \texttt{pg\_ctl start}

To start the server, waiting until the server is accepting connections:

\$ \texttt{pg\_ctl -w start}

To start the server using port 5433, and running without \texttt{fsync}, use:

\$ \texttt{pg\_ctl -o "-F -p 5433" start}

**Stopping the Server**

To stop the server, use:

\$ \texttt{pg\_ctl stop}

The \texttt{-m} option allows control over how the server shuts down:

\$ \texttt{pg\_ctl stop -m fast}

**Restarting the Server**

Restarting the server is almost equivalent to stopping the server and starting it again, except that \texttt{pg\_ctl} saves and reuses the command line options that were passed to the previously running instance. To restart the server in the simplest form, use:

\$ \texttt{pg\_ctl restart}

To restart the server, waiting for it to shut down and restart:

\$ \texttt{pg\_ctl -w restart}
To restart using port 5433, disabling fsync upon restart:

```
$ pg_ctl -o "-F -p 5433" restart
```

**Showing the Server Status**

Here is sample status output from pg_ctl:

```
$ pg_ctl status
pg_ctl: server is running (PID: 13718)
/usr/local/pgsql/bin/postgres "-D" "/usr/local/pgsql/data" "-p" "5433" "-B" "128"
```

This is the command line that would be invoked in restart mode.

**See Also**

initdb, postgres
pg_resetxlog

Name

pg_resetxlog — reset the write-ahead log and other control information of a PostgreSQL database cluster

Synopsis

pg_resetxlog [-f] [-n] [option...] [-D datadir]

Description

pg_resetxlog clears the write-ahead log (WAL) and optionally resets some other control information stored in the pg_control file. This function is sometimes needed if these files have become corrupted. It should be used only as a last resort, when the server will not start due to such corruption.

After running this command, it should be possible to start the server, but bear in mind that the database might contain inconsistent data due to partially-committed transactions. You should immediately dump your data, run initdb, and reload. After reload, check for inconsistencies and repair as needed.

This utility can only be run by the user who installed the server, because it requires read/write access to the data directory. For safety reasons, you must specify the data directory on the command line. pg_resetxlog does not use the environment variable PGDATA.

If pg_resetxlog complains that it cannot determine valid data for pg_control, you can force it to proceed anyway by specifying the -f (force) option. In this case plausible values will be substituted for the missing data. Most of the fields can be expected to match, but manual assistance might be needed for the next OID, next transaction ID and epoch, next multitransaction ID and offset, and WAL starting address fields. These fields can be set using the options discussed below. If you are not able to determine correct values for all these fields, -f can still be used, but the recovered database must be treated with even more suspicion than usual: an immediate dump and reload is imperative. Do not execute any data-modifying operations in the database before you dump, as any such action is likely to make the corruption worse.

Options

-f

Force pg_resetxlog to proceed even if it cannot determine valid data for pg_control, as explained above.

-n

The -n (no operation) option instructs pg_resetxlog to print the values reconstructed from pg_control and values about to be changed, and then exit without modifying anything. This is mainly a debugging tool, but can be useful as a sanity check before allowing pg_resetxlog to proceed for real.
-V
--version

Display version information, then exit.

-?
--help

Show help, then exit.

The following options are only needed when `pg_resetxlog` is unable to determine appropriate values by reading `pg_control`. Safe values can be determined as described below. For values that take numeric arguments, hexadecimal values can be specified by using the prefix `0x`.

-c xid,xid

Manually set the oldest and newest transaction IDs for which the commit time can be retrieved.

A safe value for the oldest transaction ID for which the commit time can be retrieved (first part) can be determined by looking for the numerically smallest file name in the directory `pg_commit_ts` under the data directory. Conversely, a safe value for the newest transaction ID for which the commit time can be retrieved (second part) can be determined by looking for the numerically greatest file name in the same directory. The file names are in hexadecimal.

-e xid_epoch

Manually set the next transaction ID’s epoch.

The transaction ID epoch is not actually stored anywhere in the database except in the field that is set by `pg_resetxlog`, so any value will work so far as the database itself is concerned. You might need to adjust this value to ensure that replication systems such as Slony-I and Skytools work correctly — if so, an appropriate value should be obtainable from the state of the downstream replicated database.

-l xlogfile

Manually set the WAL starting address.

The WAL starting address should be larger than any WAL segment file name currently existing in the directory `pg_xlog` under the data directory. These names are also in hexadecimal and have three parts. The first part is the “timeline ID” and should usually be kept the same. For example, if `0000000100000320000004A` is the largest entry in `pg_xlog`, use `-l 00000001000003200000004B` or higher.

**Note:** `pg_resetxlog` itself looks at the files in `pg_xlog` and chooses a default `-l` setting beyond the last existing file name. Therefore, manual adjustment of `-l` should only be needed if you are aware of WAL segment files that are not currently present in `pg_xlog`, such as entries in an offline archive; or if the contents of `pg_xlog` have been lost entirely.

-m mxid,mxid

Manually set the next and oldest multitransaction ID.

A safe value for the next multitransaction ID (first part) can be determined by looking for the numerically largest file name in the directory `pg_multixact/offsets` under the data directory, adding one, and then multiplying by 65536 (0x10000). Conversely, a safe value for the oldest multitransaction ID (second part of `-m`) can be determined by looking for the numerically smallest file name in the same directory and multiplying by 65536. The file names are in hexadecimal, so the easiest way to do this is to specify the option value in hexadecimal and append four zeroes.

1911
-o oid
Manually set the next OID.
There is no comparably easy way to determine a next OID that’s beyond the largest one in the
database, but fortunately it is not critical to get the next-OID setting right.

-O mxoff
Manually set the next multitransaction offset.
A safe value can be determined by looking for the numerically largest file name in the directory
pg_multixact/members under the data directory, adding one, and then multiplying by 52352
(0xCC80). The file names are in hexadecimal. There is no simple recipe such as the ones for
other options of appending zeroes.

-x xid
Manually set the next transaction ID.
A safe value can be determined by looking for the numerically largest file name in the directory
pg_clog under the data directory, adding one, and then multiplying by 1048576 (0x100000).
Note that the file names are in hexadecimal. It is usually easiest to specify the option value in
hexadecimal too. For example, if 0011 is the largest entry in pg_clog, -x 0x1200000 will
work (five trailing zeroes provide the proper multiplier).

Notes
This command must not be used when the server is running. pg_resetxlog will refuse to start up if
it finds a server lock file in the data directory. If the server crashed then a lock file might have been
left behind; in that case you can remove the lock file to allow pg_resetxlog to run. But before you
do so, make doubly certain that there is no server process still alive.

pg_resetxlog works only with servers of the same major version.

See Also
pg_controldata
**pg_rewind**

**Name**

`pg_rewind` — synchronize a PostgreSQL data directory with another data directory that was forked from it

**Synopsis**

```
pg_rewind [option...] [-D | --target-pgdata=directory] directory
|--source-pgdata=directory
|--source-server=connstr
```

**Description**

`pg_rewind` is a tool for synchronizing a PostgreSQL cluster with another copy of the same cluster, after the clusters’ timelines have diverged. A typical scenario is to bring an old master server back online after failover as a standby that follows the new master.

The result is equivalent to replacing the target data directory with the source one. Only changed blocks from relation files are copied; all other files are copied in full, including configuration files. The advantage of `pg_rewind` over taking a new base backup, or tools like rsync, is that `pg_rewind` does not require reading through unchanged blocks in the cluster. This makes it a lot faster when the database is large and only a small fraction of blocks differ between the clusters.

`pg_rewind` examines the timeline histories of the source and target clusters to determine the point where they diverged, and expects to find WAL in the target cluster’s `pg_xlog` directory reaching all the way back to the point of divergence. The point of divergence can be found either on the target timeline, the source timeline, or their common ancestor. In the typical failover scenario where the target cluster was shut down soon after the divergence, this is not a problem, but if the target cluster ran for a long time after the divergence, the old WAL files might no longer be present. In that case, they can be manually copied from the WAL archive to the `pg_xlog` directory, or fetched on startup by configuring `recovery.conf`. The use of `pg_rewind` is not limited to failover, e.g. a standby server can be promoted, run some write transactions, and then rewinded to become a standby again.

When the target server is started for the first time after running `pg_rewind`, it will go into recovery mode and replay all WAL generated in the source server after the point of divergence. If some of the WAL was no longer available in the source server when `pg_rewind` was run, and therefore could not be copied by the `pg_rewind` session, it must be made available when the target server is started. This can be done by creating a `recovery.conf` file in the target data directory with a suitable `restore_command`.

`pg_rewind` requires that the target server either has the `wal_log_hints` option enabled in `postgresql.conf` or data checksums enabled when the cluster was initialized with `initdb`. Neither of these are currently on by default. `full_page_writes` must also be set to `on`, but is enabled by default.
Warning

If `pg_rewind` fails while processing, then the data folder of the target is likely not in a state that can be recovered. In such a case, taking a new fresh backup is recommended.

`pg_rewind` will fail immediately if it finds files it cannot write directly to. This can happen for example when the source and the target server use the same file mapping for read-only SSL keys and certificates. If such files are present on the target server it is recommended to remove them before running `pg_rewind`. After doing the rewind, some of those files may have been copied from the source, in which case it may be necessary to remove the data copied and restore back the set of links used before the rewind.

Options

`pg_rewind` accepts the following command-line arguments:

```
-D directory
--target-pgdata=directory
```

This option specifies the target data directory that is synchronized with the source. The target server must be shut down cleanly before running `pg_rewind`.

```
--source-pgdata=directory
```

Specifies the file system path to the data directory of the source server to synchronize the target with. This option requires the source server to be cleanly shut down.

```
--source-server=connstr
```

Specifies a libpq connection string to connect to the source PostgreSQL server to synchronize the target with. The connection must be a normal (non-replication) connection with superuser access. This option requires the source server to be running and not in recovery mode.

```
-n
--dry-run
```

Do everything except actually modifying the target directory.

```
-P
--progress
```

Enables progress reporting. Turning this on will deliver an approximate progress report while copying data from the source cluster.

```
--debug
```

Print verbose debugging output that is mostly useful for developers debugging `pg_rewind`.

```
-V
--version
```

Display version information, then exit.

```
-?
--help
```

Show help, then exit.
Environment

When --source-server option is used, pg_rewind also uses the environment variables supported by libpq (see Section 32.14).

Notes

When executing pg_rewind using an online cluster as source which has been recently promoted, it is necessary to execute a CHECKPOINT after promotion so as its control file reflects up-to-date timeline information, which is used by pg_rewind to check if the target cluster can be rewound using the designated source cluster.

How it works

The basic idea is to copy all file system-level changes from the source cluster to the target cluster:

1. Scan the WAL log of the target cluster, starting from the last checkpoint before the point where the source cluster’s timeline history forked off from the target cluster. For each WAL record, record each data block that was touched. This yields a list of all the data blocks that were changed in the target cluster, after the source cluster forked off.

2. Copy all those changed blocks from the source cluster to the target cluster, either using direct file system access (--source-pgdata) or SQL (--source-server).

3. Copy all other files such as pg_clog and configuration files from the source cluster to the target cluster (everything except the relation files).

4. Apply the WAL from the source cluster, starting from the checkpoint created at failover. (Strictly speaking, pg_rewind doesn’t apply the WAL, it just creates a backup label file that makes PostgreSQL start by replaying all WAL from that checkpoint forward.)
pg_test_fsync

Name

pg_test_fsync — determine fastest wal_sync_method for PostgreSQL

Synopsis

pg_test_fsync [option...]

Description

pg_test_fsync is intended to give you a reasonable idea of what the fastest wal_sync_method is on your specific system, as well as supplying diagnostic information in the event of an identified I/O problem. However, differences shown by pg_test_fsync might not make any significant difference in real database throughput, especially since many database servers are not speed-limited by their transaction logs. pg_test_fsync reports average file sync operation time in microseconds for each wal_sync_method, which can also be used to inform efforts to optimize the value of commit_delay.

Options

pg_test_fsync accepts the following command-line options:

-\f
--filename

   Specifies the file name to write test data in. This file should be in the same file system that the pg_xlog directory is or will be placed in. (pg_xlog contains the WAL files.) The default is pg_test_fsync.out in the current directory.

-\s
--secs-per-test

   Specifies the number of seconds for each test. The more time per test, the greater the test’s accuracy, but the longer it takes to run. The default is 5 seconds, which allows the program to complete in under 2 minutes.

-\V
--version

   Print the pg_test_fsync version and exit.

-\?
--help

   Show help about pg_test_fsync command line arguments, and exit.
See Also

postgres
pg_test_timing

Name
pg_test_timing — measure timing overhead

Synopsis
pg_test_timing[option...]

Description
pg_test_timing is a tool to measure the timing overhead on your system and confirm that the system
time never moves backwards. Systems that are slow to collect timing data can give less accurate
EXPLAIN ANALYZE results.

Options
pg_test_timing accepts the following command-line options:

-d duration
--duration=duration
  Specifies the test duration, in seconds. Longer durations give slightly better accuracy, and are
  more likely to discover problems with the system clock moving backwards. The default test
duration is 3 seconds.

-V
--version
  Print the pg_test_timing version and exit.

-?
--help
  Show help about pg_test_timing command line arguments, and exit.

Usage

Interpreting results
Good results will show most (>90%) individual timing calls take less than one microsecond. Average
per loop overhead will be even lower, below 100 nanoseconds. This example from an Intel i7-860
system using a TSC clock source shows excellent performance:

Testing timing overhead for 3 seconds.
Per loop time including overhead: 35.96 nsec
Histogram of timing durations:
Note that different units are used for the per loop time than the histogram. The loop can have resolution within a few nanoseconds (nsec), while the individual timing calls can only resolve down to one microsecond (usec).

### Measuring executor timing overhead

When the query executor is running a statement using `EXPLAIN ANALYZE`, individual operations are timed as well as showing a summary. The overhead of your system can be checked by counting rows with the `psql` program:

```sql
CREATE TABLE t AS SELECT * FROM generate_series(1,100000);
\timing
SELECT COUNT(*) FROM t;
EXPLAIN ANALYZE SELECT COUNT(*) FROM t;
```

The i7-860 system measured runs the count query in 9.8 ms while the `EXPLAIN ANALYZE` version takes 16.6 ms, each processing just over 100,000 rows. That 6.8 ms difference means the timing overhead per row is 68 ns, about twice what `pg_test_timing` estimated it would be. Even that relatively small amount of overhead is making the fully timed count statement take almost 70% longer. On more substantial queries, the timing overhead would be less problematic.

### Changing time sources

On some newer Linux systems, it’s possible to change the clock source used to collect timing data at any time. A second example shows the slowdown possible from switching to the slower acpi_pm time source, on the same system used for the fast results above:

```bash
# cat /sys/devices/system/clocksource/clocksource0/available_clocksource
tsc hpet acpi_pm
# echo acpi_pm > /sys/devices/system/clocksource/clocksource0/current_clocksource
# pg_test_timing
Per loop time including overhead: 722.92 nsec
Histogram of timing durations:
< usec % of total count
 1 27.84870 1155682
 2 72.05956 2990371
 4 0.07810 3241
 8 0.01357 563
16 0.00007 3
```

In this configuration, the sample `EXPLAIN ANALYZE` above takes 115.9 ms. That’s 1061 nsec of timing overhead, again a small multiple of what’s measured directly by this utility. That much timing
overhead means the actual query itself is only taking a tiny fraction of the accounted for time, most of it is being consumed in overhead instead. In this configuration, any EXPLAIN ANALYZE totals involving many timed operations would be inflated significantly by timing overhead.

FreeBSD also allows changing the time source on the fly, and it logs information about the timer selected during boot:

```
# dmesg | grep "Timecounter"
Timecounter "ACPI-fast" frequency 3579545 Hz quality 900
Timecounter "i8254" frequency 1193182 Hz quality 0
Timecounters tick every 10.000 msec
Timecounter "TSC" frequency 2531787134 Hz quality 800
# sysctl kern.timecounter.hardware=TSC
kern.timecounter.hardware: ACPI-fast -> TSC
```

Other systems may only allow setting the time source on boot. On older Linux systems the "clock" kernel setting is the only way to make this sort of change. And even on some more recent ones, the only option you’ll see for a clock source is "jiffies". Jiffies are the older Linux software clock implementation, which can have good resolution when it’s backed by fast enough timing hardware, as in this example:

```
$ cat /sys/devices/system/clocksource/clocksource0/available_clocksource
jiffies
$ dmesg | grep time.c
time.c: Using 3.579545 MHz WALL PM GTOD PIT/TSC timer.
time.c: Detected 2400.153 MHz processor.
$ pg_test_timing
Testing timing overhead for 3 seconds.
Per timing duration including loop overhead: 97.75 ns
Histogram of timing durations:
< usec % of total count
 1  90.23734  27694571
 2  9.75277  2993204
 4  0.00981  3010
 8  0.00007  22
16  0.00000  1
32  0.00000  1
```

**Clock hardware and timing accuracy**

Collecting accurate timing information is normally done on computers using hardware clocks with various levels of accuracy. With some hardware the operating systems can pass the system clock time almost directly to programs. A system clock can also be derived from a chip that simply provides timing interrupts, periodic ticks at some known time interval. In either case, operating system kernels provide a clock source that hides these details. But the accuracy of that clock source and how quickly it can return results varies based on the underlying hardware.

Inaccurate time keeping can result in system instability. Test any change to the clock source very carefully. Operating system defaults are sometimes made to favor reliability over best accuracy. And if you are using a virtual machine, look into the recommended time sources compatible with it. Virtual hardware faces additional difficulties when emulating timers, and there are often per operating system settings suggested by vendors.
The Time Stamp Counter (TSC) clock source is the most accurate one available on current generation CPUs. It’s the preferred way to track the system time when it’s supported by the operating system and the TSC clock is reliable. There are several ways that TSC can fail to provide an accurate timing source, making it unreliable. Older systems can have a TSC clock that varies based on the CPU temperature, making it unusable for timing. Trying to use TSC on some older multicore CPUs can give a reported time that’s inconsistent among multiple cores. This can result in the time going backwards, a problem this program checks for. And even the newest systems can fail to provide accurate TSC timing with very aggressive power saving configurations.

Newer operating systems may check for the known TSC problems and switch to a slower, more stable clock source when they are seen. If your system supports TSC time but doesn’t default to that, it may be disabled for a good reason. And some operating systems may not detect all the possible problems correctly, or will allow using TSC even in situations where it’s known to be inaccurate.

The High Precision Event Timer (HPET) is the preferred timer on systems where it’s available and TSC is not accurate. The timer chip itself is programmable to allow up to 100 nanosecond resolution, but you may not see that much accuracy in your system clock.

Advanced Configuration and Power Interface (ACPI) provides a Power Management (PM) Timer, which Linux refers to as the acpi_pm. The clock derived from acpi_pm will at best provide 300 nanosecond resolution.

Timers used on older PC hardware include the 8254 Programmable Interval Timer (PIT), the real-time clock (RTC), the Advanced Programmable Interrupt Controller (APIC) timer, and the Cyclone timer. These timers aim for millisecond resolution.

See Also

EXPLAIN
pg_upgrade

Name
pg_upgrade — upgrade a PostgreSQL server instance

Synopsis
pg_upgrade -b oldbindir -B newbindir -d olddatadir -D newdatadir [option...]

Description
pg_upgrade (formerly called pg_migrator) allows data stored in PostgreSQL data files to be upgraded to a later PostgreSQL major version without the data dump/reload typically required for major version upgrades, e.g. from 8.4.7 to the current major release of PostgreSQL. It is not required for minor version upgrades, e.g. from 9.0.1 to 9.0.4.

Major PostgreSQL releases regularly add new features that often change the layout of the system tables, but the internal data storage format rarely changes. pg_upgrade uses this fact to perform rapid upgrades by creating new system tables and simply reusing the old user data files. If a future major release ever changes the data storage format in a way that makes the old data format unreadable, pg_upgrade will not be usable for such upgrades. (The community will attempt to avoid such situations.)

pg_upgrade does its best to make sure the old and new clusters are binary-compatible, e.g. by checking for compatible compile-time settings, including 32/64-bit binaries. It is important that any external modules are also binary compatible, though this cannot be checked by pg_upgrade.

pg_upgrade supports upgrades from 8.4.X and later to the current major release of PostgreSQL, including snapshot and alpha releases.

Options
pg_upgrade accepts the following command-line arguments:

- b bindir
  --old-bindir=bindir
      the old PostgreSQL executable directory; environment variable PGINOLD

- B bindir
  --new-bindir=bindir
      the new PostgreSQL executable directory; environment variable PGINNEW

- c
  --check
      check clusters only, don’t change any data
pg_upgrade

-d datadir
--old-datadir=datadir
  the old cluster data directory; environment variable PGDATAOLD

-D datadir
--new-datadir=datadir
  the new cluster data directory; environment variable PGDATANEW

-j
--jobs
  number of simultaneous processes or threads to use

-k
--link
  use hard links instead of copying files to the new cluster

-o options
--old-options options
  options to be passed directly to the old postgres command; multiple option invocations are appended

-O options
--new-options options
  options to be passed directly to the new postgres command; multiple option invocations are appended

-p port
--old-port=port
  the old cluster port number; environment variable PGPORTOLD

-P port
--new-port=port
  the new cluster port number; environment variable PGPORTNEW

-r
--retain
  retain SQL and log files even after successful completion

-U username
--username=username
  cluster's install user name; environment variable PGUSER

-v
--verbose
  enable verbose internal logging

-V
--version
  display version information, then exit

-?
--help
  show help, then exit
Usage

These are the steps to perform an upgrade with pg_upgrade:

1. Optionally move the old cluster

   If you are using a version-specific installation directory, e.g. /opt/PostgreSQL/9.1, you do not need to move the old cluster. The graphical installers all use version-specific installation directories.

   If your installation directory is not version-specific, e.g. /usr/local/pgsql, it is necessary to move the current PostgreSQL install directory so it does not interfere with the new PostgreSQL installation. Once the current PostgreSQL server is shut down, it is safe to rename the PostgreSQL installation directory; assuming the old directory is /usr/local/pgsql, you can do:

   mv /usr/local/pgsql /usr/local/pgsql.old

   to rename the directory.

2. For source installs, build the new version

   Build the new PostgreSQL source with configure flags that are compatible with the old cluster. pg_upgrade will check pg_controldata to make sure all settings are compatible before starting the upgrade.

3. Install the new PostgreSQL binaries

   Install the new server’s binaries and support files. pg_upgrade is included in a default installation.

   For source installs, if you wish to install the new server in a custom location, use the prefix variable:

   make prefix=/usr/local/pgsql.new install

4. Initialize the new PostgreSQL cluster

   Initialize the new cluster using initdb. Again, use compatible initdb flags that match the old cluster. Many prebuilt installers do this step automatically. There is no need to start the new cluster.

5. Install custom shared object files

   Install any custom shared object files (or DLLs) used by the old cluster into the new cluster, e.g. pgcrypto.so, whether they are from contrib or some other source. Do not install the schema definitions, e.g. CREATE EXTENSION pgcrypto, because these will be upgraded from the old cluster. Also, any custom full text search files (dictionary, synonym, thesaurus, stop words) must also be copied to the new cluster.

6. Adjust authentication

   pg_upgrade will connect to the old and new servers several times, so you might want to set authentication to peer in pg_hba.conf or use a ~/.pgpass file (see Section 32.15).

7. Stop both servers

   Make sure both database servers are stopped using, on Unix, e.g.:

   pg_ctl -D /opt/PostgreSQL/8.4 stop
   pg_ctl -D /opt/PostgreSQL/9.0 stop

   or on Windows, using the proper service names:

   NET STOP postgresql-8.4
   NET STOP postgresql-9.0

   Streaming replication and log-shipping standby servers can remain running until a later step.
8. Prepare for standby server upgrades

If you are upgrading standby servers using methods outlined in section step 10, verify that the old standby servers are caught up by running pg_controldata against the old primary and standby clusters. Verify that the “Latest checkpoint location” values match in all clusters. (There will be a mismatch if old standby servers were shut down before the old primary or if the old standby servers are still running.) Also, change wal_level to replica in the postgresql.conf file on the new primary cluster.

9. Run pg_upgrade

Always run the pg_upgrade binary of the new server, not the old one. pg_upgrade requires the specification of the old and new cluster’s data and executable (bin) directories. You can also specify user and port values, and whether you want the data linked instead of copied (the default).

If you use link mode, the upgrade will be much faster (no file copying) and use less disk space, but you will not be able to access your old cluster once you start the new cluster after the upgrade. Link mode also requires that the old and new cluster data directories be in the same file system. (Tablespaces and pg_xlog can be on different file systems.) See pg_upgrade --help for a full list of options.

The --jobs option allows multiple CPU cores to be used for copying/linking of files and to dump and reload database schemas in parallel; a good place to start is the maximum of the number of CPU cores and tablespaces. This option can dramatically reduce the time to upgrade a multi-database server running on a multiprocessor machine.

For Windows users, you must be logged into an administrative account, and then start a shell as the postgres user and set the proper path:

```bash
RUNAS /USER:postgres "CMD.EXE"
SET PATH=%PATH%;C:\Program Files\PostgreSQL\9.0\bin;
```

and then run pg_upgrade with quoted directories, e.g.:

```bash
pg_upgrade.exe
   --old-datadir "C:/Program Files/PostgreSQL/8.4/data"
   --new-datadir "C:/Program Files/PostgreSQL/9.0/data"
   --old-bindir "C:/Program Files/PostgreSQL/8.4/bin"
   --new-bindir "C:/Program Files/PostgreSQL/9.0/bin"
```

Once started, pg_upgrade will verify the two clusters are compatible and then do the upgrade. You can use pg_upgrade --check to perform only the checks, even if the old server is still running. pg_upgrade --check will also outline any manual adjustments you will need to make after the upgrade. If you are going to be using link mode, you should use the --link option with --check to enable link-mode-specific checks. pg_upgrade requires write permission in the current directory.

Obviously, no one should be accessing the clusters during the upgrade. pg_upgrade defaults to running servers on port 50432 to avoid unintended client connections. You can use the same port number for both clusters when doing an upgrade because the old and new clusters will not be running at the same time. However, when checking an old running server, the old and new port numbers must be different.

If an error occurs while restoring the database schema, pg_upgrade will exit and you will have to revert to the old cluster as outlined in step 16 below. To try pg_upgrade again, you will need to modify the old cluster so the pg_upgrade schema restore succeeds. If the problem is a contrib module, you might need to uninstall the contrib module from the old cluster and install it in the new cluster after the upgrade, assuming the module is not being used to store user data.

10. Upgrade Streaming Replication and Log-Shipping standby servers
If you used link mode and have Streaming Replication (see Section 26.2.5) or Log-Shipping (see Section 26.2) standby servers, you can follow these steps to quickly upgrade them. You will not be running pg_upgrade on the standby servers, but rather rsync on the primary. Do not start any servers yet.

If you did not use link mode, do not have or do not want to use rsync, or want an easier solution, skip the instructions in this section and simply recreate the standby servers once pg_upgrade completes and the new primary is running.

   a. Install the new PostgreSQL binaries on standby servers
      Make sure the new binaries and support files are installed on all standby servers.
   
   b. Make sure the new standby data directories do not exist
      Make sure the new standby data directories do not exist or are empty. If initdb was run, delete the standby servers’ new data directories.
   
   c. Install custom shared object files
      Install the same custom shared object files on the new standbys that you installed in the new primary cluster.
   
   d. Stop standby servers
      If the standby servers are still running, stop them now using the above instructions.
   
   e. Save configuration files
      Save any configuration files from the old standbys’ data directories you need to keep, e.g. postgresql.conf, recovery.conf, because these will be overwritten or removed in the next step.
   
   f. Run rsync
      When using link mode, standby servers can be quickly upgraded using rsync. To accomplish this, from a directory on the primary server that is above the old and new database cluster directories, run this on the primary for each standby server:
      
      rsync --archive --delete --hard-links --size-only --no-inc-recursive old_cluster new_cluster remote_dir
      
      where old_cluster and new_cluster are relative to the current directory on the primary, and remote_dir is above the old and new cluster directories on the standby. The directory structure under the specified directories on the primary and standbys must match. Consult the rsync manual page for details on specifying the remote directory, e.g.
      
      rsync --archive --delete --hard-links --size-only --no-inc-recursive /opt/PostgreSQL/9.6 /opt/PostgreSQL/9.6 standby.example.com:/opt/PostgreSQL
      
      You can verify what the command will do using rsync’s --dry-run option. While rsync must be run on the primary for at least one standby, it is possible to run rsync on an upgraded standby to upgrade other standbys, as long as the upgraded standby has not been started.
      
      What this does is to record the links created by pg_upgrade’s link mode that connect files in the old and new clusters on the primary server. It then finds matching files in the standby’s old cluster and creates links for them in the standby’s new cluster. Files that were not linked on the primary are copied from the primary to the standby. (They are usually small.) This provides rapid standby upgrades. Unfortunately, rsync needlessly copies files associated with temporary and unlogged tables because these files don’t normally exist on standby servers.
      
      If you have tablespaces, you will need to run a similar rsync command for each tablespace directory, e.g.:
rsync --archive --delete --hard-links --size-only --no-inc-recursive /vol1/pg_tblsp/PG_9.5_201510051 /vol1/pg_tblsp/PG_9.6_201608131 standby.example.com:/vol1/pg_tblsp

If you have relocated pg_xlog outside the data directories, rsync must be run on those directories too.

g. Configure streaming replication and log-shipping standby servers

Configure the servers for log shipping. (You do not need to run pg_start_backup() and pg_stop_backup() or take a file system backup as the standbys are still synchronized with the primary.)

11. Restore pg_hba.conf

If you modified pg_hba.conf, restore its original settings. It might also be necessary to adjust other configuration files in the new cluster to match the old cluster, e.g. postgresql.conf.

12. Start the new server

The new server can now be safely started, and then any rsync’ed standby servers.

13. Post-Upgrade processing

If any post-upgrade processing is required, pg_upgrade will issue warnings as it completes. It will also generate script files that must be run by the administrator. The script files will connect to each database that needs post-upgrade processing. Each script should be run using:

```
psql --username postgres --file script.sql postgres
```

The scripts can be run in any order and can be deleted once they have been run.

---

**Caution**

In general it is unsafe to access tables referenced in rebuild scripts until the rebuild scripts have run to completion; doing so could yield incorrect results or poor performance. Tables not referenced in rebuild scripts can be accessed immediately.

---

14. Statistics

Because optimizer statistics are not transferred by pg_upgrade, you will be instructed to run a command to regenerate that information at the end of the upgrade. You might need to set connection parameters to match your new cluster.

15. Delete old cluster

Once you are satisfied with the upgrade, you can delete the old cluster’s data directories by running the script mentioned when pg_upgrade completes. (Automatic deletion is not possible if you have user-defined tablespaces inside the old data directory.) You can also delete the old installation directories (e.g. bin, share).

16. Reverting to old cluster

If, after running pg_upgrade, you wish to revert to the old cluster, there are several options:

- If the --check option was used, the old cluster was unmodified; it can be restarted.
- If the --link option was not used, the old cluster was unmodified; it can be restarted.
- If the --link option was used, the data files might be shared between the old and new cluster:
  - If pg_upgrade aborted before linking started, the old cluster was unmodified; it can be restarted.
  - If you did not start the new cluster, the old cluster was unmodified except that, when linking started, a .old suffix was appended to $PGDATA/global/pg_control. To reuse the old
cluster, remove the .old suffix from $PGDATA/global/pg_control; you can then restart
the old cluster.

- If you did start the new cluster, it has written to shared files and it is unsafe to use the old
cluster. The old cluster will need to be restored from backup in this case.

Notes

pg_upgrade does not support upgrading of databases containing these reg* OID-referencing
system data types: regproc, regprocedure, regoper, regoperator, regconfig, and
regdictionary. (regtype can be upgraded.)

All failure, rebuild, and reindex cases will be reported by pg_upgrade if they affect your installation;
post-upgrade scripts to rebuild tables and indexes will be generated automatically. If you are trying to
automate the upgrade of many clusters, you should find that clusters with identical database schemas
require the same post-upgrade steps for all cluster upgrades; this is because the post-upgrade steps are
based on the database schemas, and not user data.

For deployment testing, create a schema-only copy of the old cluster, insert dummy data, and upgrade
that.

If you are upgrading a pre-PostgreSQL 9.2 cluster that uses a configuration-file-only directory, you
must pass the real data directory location to pg_upgrade, and pass the configuration directory location
to the server, e.g. -d /real-data-directory -o '-D /configuration-directory'.

If using a pre-9.1 old server that is using a non-default Unix-domain socket directory or a default that
differs from the default of the new cluster, set PGHOST to point to the old server’s socket location.
(This is not relevant on Windows.)

If you want to use link mode and you do not want your old cluster to be modified when the new cluster
is started, make a copy of the old cluster and upgrade that in link mode. To make a valid copy of the
old cluster, use rsync to create a dirty copy of the old cluster while the server is running, then shut
down the old server and run rsync --checksum again to update the copy with any changes to make
it consistent. (--checksum is necessary because rsync only has file modification-time granularity of
one second.) You might want to exclude some files, e.g. postmaster.pid, as documented in Section
25.3.3. If your file system supports file system snapshots or copy-on-write file copies, you can use that
to make a backup of the old cluster and tablespaces, though the snapshot and copies must be created
simultaneously or while the database server is down.

See Also

initdb, pg_ctl, pg_dump, postgres
**pg_xlogdump**

**Name**

pg_xlogdump — display a human-readable rendering of the write-ahead log of a PostgreSQL database cluster

**Synopsis**

pg_xlogdump [option...][startseg [endseg]]

**Description**

pg_xlogdump displays the write-ahead log (WAL) and is mainly useful for debugging or educational purposes.

This utility can only be run by the user who installed the server, because it requires read-only access to the data directory.

**Options**

The following command-line options control the location and format of the output:

- `startseg`
  - Start reading at the specified log segment file. This implicitly determines the path in which files will be searched for, and the timeline to use.

- `endseg`
  - Stop after reading the specified log segment file.

- `-b`  
  - `--bkp-details`
    - Output detailed information about backup blocks.

- `-e end`  
  - `--end=end`
    - Stop reading at the specified log position, instead of reading to the end of the log stream.

- `-f`  
  - `--follow`
    - After reaching the end of valid WAL, keep polling once per second for new WAL to appear.

- `-n limit`  
  - `--limit=limit`
    - Display the specified number of records, then stop.
pg_xlogdump

-p path
--path=path

Specifies a directory to search for log segment files or a directory with a pg_xlog subdirectory that contains such files. The default is to search in the current directory, the pg_xlog subdirectory of the current directory, and the pg_xlog subdirectory of PGDATA.

-r rmgr
--rmgr=rmgr

Only display records generated by the specified resource manager. If list is passed as name, print a list of valid resource manager names, and exit.

-s start
--start=start

Log position at which to start reading. The default is to start reading the first valid log record found in the earliest file found.

-t timeline
--timeline=timeline

Timeline from which to read log records. The default is to use the value in startseg, if that is specified; otherwise, the default is 1.

-V
--version

Print the pg_xlogdump version and exit.

-x xid
--xid=xid

Only display records marked with the given transaction ID.

-z
--stats[=record]

Display summary statistics (number and size of records and full-page images) instead of individual records. Optionally generate statistics per-record instead of per-rmgr.

-?
--help

Show help about pg_xlogdump command line arguments, and exit.

Notes

Can give wrong results when the server is running.

Only the specified timeline is displayed (or the default, if none is specified). Records in other timelines are ignored.

pg_xlogdump cannot read WAL files with suffix .partial. If those files need to be read, .partial suffix needs to be removed from the file name.
See Also
Section 30.5
**postgres**

**Name**

*postgres* — PostgreSQL database server

**Synopsis**

```
postgres [option...]
```

**Description**

*postgres* is the PostgreSQL database server. In order for a client application to access a database it connects (over a network or locally) to a running *postgres* instance. The *postgres* instance then starts a separate server process to handle the connection.

One *postgres* instance always manages the data of exactly one database cluster. A database cluster is a collection of databases that is stored at a common file system location (the “data area”). More than one *postgres* instance can run on a system at one time, so long as they use different data areas and different communication ports (see below). When *postgres* starts it needs to know the location of the data area. The location must be specified by the `-D` option or the `PGDATA` environment variable; there is no default. Typically, `-D` or `PGDATA` points directly to the data area directory created by initdb. Other possible file layouts are discussed in Section 19.2.

By default *postgres* starts in the foreground and prints log messages to the standard error stream. In practical applications *postgres* should be started as a background process, perhaps at boot time.

The *postgres* command can also be called in single-user mode. The primary use for this mode is during bootstrapping by initdb. Sometimes it is used for debugging or disaster recovery; note that running a single-user server is not truly suitable for debugging the server, since no realistic interprocess communication and locking will happen. When invoked in single-user mode from the shell, the user can enter queries and the results will be printed to the screen, but in a form that is more useful for developers than end users. In the single-user mode, the session user will be set to the user with ID 1, and implicit superuser powers are granted to this user. This user does not actually have to exist, so the single-user mode can be used to manually recover from certain kinds of accidental damage to the system catalogs.

**Options**

*postgres* accepts the following command-line arguments. For a detailed discussion of the options consult Chapter 19. You can save typing most of these options by setting up a configuration file. Some (safe) options can also be set from the connecting client in an application-dependent way to apply only for that session. For example, if the environment variable `PGOPTIONS` is set, then libpq-based clients will pass that string to the server, which will interpret it as *postgres* command-line options.
General Purpose

-B nbuffers
Sets the number of shared buffers for use by the server processes. The default value of this parameter is chosen automatically by initdb. Specifying this option is equivalent to setting the shared_buffers configuration parameter.

-c name=value
Sets a named run-time parameter. The configuration parameters supported by PostgreSQL are described in Chapter 19. Most of the other command line options are in fact short forms of such a parameter assignment. -c can appear multiple times to set multiple parameters.

-C name
Prints the value of the named run-time parameter, and exits. (See the -c option above for details.) This can be used on a running server, and returns values from postgresql.conf, modified by any parameters supplied in this invocation. It does not reflect parameters supplied when the cluster was started.
This option is meant for other programs that interact with a server instance, such as pg_ctl, to query configuration parameter values. User-facing applications should instead use SHOW or the pg_settings view.

-d debug-level
Sets the debug level. The higher this value is set, the more debugging output is written to the server log. Values are from 1 to 5. It is also possible to pass -d 0 for a specific session, which will prevent the server log level of the parent postgres process from being propagated to this session.

-D datadir
Specifies the file system location of the database configuration files. See Section 19.2 for details.

-e
Sets the default date style to “European”, that is DMY ordering of input date fields. This also causes the day to be printed before the month in certain date output formats. See Section 8.5 for more information.

-F
Disables fsync calls for improved performance, at the risk of data corruption in the event of a system crash. Specifying this option is equivalent to disabling the fsync configuration parameter. Read the detailed documentation before using this!

-h hostname
Specifies the IP host name or address on which postgres is to listen for TCP/IP connections from client applications. The value can also be a comma-separated list of addresses, or * to specify listening on all available interfaces. An empty value specifies not listening on any IP addresses, in which case only Unix-domain sockets can be used to connect to the server. Defaults to listening only on localhost. Specifying this option is equivalent to setting the listen_addresses configuration parameter.

-i
Allows remote clients to connect via TCP/IP (Internet domain) connections. Without this option, only local connections are accepted. This option is equivalent to setting listen_addresses to * in postgresql.conf or via -h.
This option is deprecated since it does not allow access to the full functionality of
listen_addresses. It’s usually better to set listen_addresses directly.

-\texttt{k directory}

Specifies the directory of the Unix-domain socket on which \texttt{postgres} is to listen for connections from client applications. The value can also be a comma-separated list of directories. An empty value specifies not listening on any Unix-domain sockets, in which case only TCP/IP sockets can be used to connect to the server. The default value is normally /tmp, but that can be changed at build time. Specifying this option is equivalent to setting the unix_socket_directories configuration parameter.

-\texttt{l}

Enables secure connections using SSL. PostgreSQL must have been compiled with support for SSL for this option to be available. For more information on using SSL, refer to Section 18.9.

-\texttt{N max-connections}

Sets the maximum number of client connections that this server will accept. The default value of this parameter is chosen automatically by initdb. Specifying this option is equivalent to setting the max_connections configuration parameter.

-\texttt{o extra-options}

The command-line-style arguments specified in extra-options are passed to all server processes started by this \texttt{postgres} process.

Spaces within extra-options are considered to separate arguments, unless escaped with a backslash (\); write \texttt{\textbackslash} to represent a literal backslash. Multiple arguments can also be specified via multiple uses of \texttt{-o}.

The use of this option is obsolete; all command-line options for server processes can be specified directly on the \texttt{postgres} command line.

-\texttt{p port}

Specifies the TCP/IP port or local Unix domain socket file extension on which \texttt{postgres} is to listen for connections from client applications. Defaults to the value of the PGPORT environment variable, or if PGPORT is not set, then defaults to the value established during compilation (normally 5432). If you specify a port other than the default port, then all client applications must specify the same port using either command-line options or PGPORT.

-\texttt{s}

Print time information and other statistics at the end of each command. This is useful for benchmarking or for use in tuning the number of buffers.

-\texttt{S work-mem}

Specifies the amount of memory to be used by internal sorts and hashes before resorting to temporary disk files. See the description of the work_mem configuration parameter in Section 19.4.1.

-\texttt{V}

-\texttt{--version}

Print the postgres version and exit.

-\texttt{--name=value}

Sets a named run-time parameter; a shorter form of \texttt{-c}. 

1934
This option dumps out the server’s internal configuration variables, descriptions, and defaults in tab-delimited COPY format. It is designed primarily for use by administration tools.

Show help about postgres command line arguments, and exit.

Semi-internal Options

The options described here are used mainly for debugging purposes, and in some cases to assist with recovery of severely damaged databases. There should be no reason to use them in a production database setup. They are listed here only for use by PostgreSQL system developers. Furthermore, these options might change or be removed in a future release without notice.

Forbids the use of particular scan and join methods: s and i disable sequential and index scans respectively, o, b and t disable index-only scans, bitmap index scans, and TID scans respectively, while n, m, and h disable nested-loop, merge and hash joins respectively.

Neither sequential scans nor nested-loop joins can be disabled completely; the -fs and -fn options simply discourage the optimizer from using those plan types if it has any other alternative.

This option is for debugging problems that cause a server process to die abnormally. The ordinary strategy in this situation is to notify all other server processes that they must terminate and then reinitialize the shared memory and semaphores. This is because an errant server process could have corrupted some shared state before terminating. This option specifies that postgres will not reinitialize shared data structures. A knowledgeable system programmer can then use a debugger to examine shared memory and semaphore state.

Allows the structure of system tables to be modified. This is used by initdb.

Ignore system indexes when reading system tables, but still update the indexes when modifying the tables. This is useful when recovering from damaged system indexes.

Print timing statistics for each query relating to each of the major system modules. This option cannot be used together with the -s option.

This option is for debugging problems that cause a server process to die abnormally. The ordinary strategy in this situation is to notify all other server processes that they must terminate and then reinitialize the shared memory and semaphores. This is because an errant server process could have corrupted some shared state before terminating. This option specifies that postgres will stop all other server processes by sending the signal SIGSTOP, but will not cause them to terminate. This permits system programmers to collect core dumps from all server processes by hand.
**Options for Single-User Mode**

The following options only apply to the single-user mode (see *Single-User Mode*).

---

**--single**

Selects the single-user mode. This must be the first argument on the command line.

**database**

Specifies the name of the database to be accessed. This must be the last argument on the command line. If it is omitted it defaults to the user name.

**-E**

Echo all commands to standard output before executing them.

**-j**

Use semicolon followed by two newlines, rather than just newline, as the command entry terminator.

**-r filename**

Send all server log output to `filename`. This option is only honored when supplied as a command-line option.

---

**Environment**

**PGCLIENTENCODING**

Default character encoding used by clients. (The clients can override this individually.) This value can also be set in the configuration file.

**PGDATA**

Default data directory location

**PGDATESTYLE**

Default value of the DateStyle run-time parameter. (The use of this environment variable is deprecated.)

**PGPORT**

Default port number (preferably set in the configuration file)
Diagnostics

A failure message mentioning `semget` or `shmget` probably indicates you need to configure your kernel to provide adequate shared memory and semaphores. For more discussion see Section 18.4. You might be able to postpone reconfiguring your kernel by decreasing `shared_buffers` to reduce the shared memory consumption of PostgreSQL, and/or by reducing `max_connections` to reduce the semaphore consumption.

A failure message suggesting that another server is already running should be checked carefully, for example by using the command

```
$ ps ax | grep postgres
```

or

```
$ ps -ef | grep postgres
```

depending on your system. If you are certain that no conflicting server is running, you can remove the lock file mentioned in the message and try again.

A failure message indicating inability to bind to a port might indicate that that port is already in use by some non-PostgreSQL process. You might also get this error if you terminate `postgres` and immediately restart it using the same port; in this case, you must simply wait a few seconds until the operating system closes the port before trying again. Finally, you might get this error if you specify a port number that your operating system considers to be reserved. For example, many versions of Unix consider port numbers under 1024 to be “trusted” and only permit the Unix superuser to access them.

Notes

The utility command `pg_ctl` can be used to start and shut down the `postgres` server safely and comfortably.

If at all possible, do not use `SIGKILL` to kill the main `postgres` server. Doing so will prevent `postgres` from freeing the system resources (e.g., shared memory and semaphores) that it holds before terminating. This might cause problems for starting a fresh `postgres` run.

To terminate the `postgres` server normally, the signals `SIGTERM`, `SIGINT`, or `SIGQUIT` can be used. The first will wait for all clients to terminate before quitting, the second will forcefully disconnect all clients, and the third will quit immediately without proper shutdown, resulting in a recovery run during restart.

The `SIGHUP` signal will reload the server configuration files. It is also possible to send `SIGHUP` to an individual server process, but that is usually not sensible.

To cancel a running query, send the `SIGINT` signal to the process running that command. To terminate a backend process cleanly, send `SIGTERM` to that process. See also `pg_cancel_backend` and `pg_terminate_backend` in Section 9.26.2 for the SQL-callable equivalents of these two actions.

The `postgres` server uses `SIGQUIT` to tell subordinate server processes to terminate without normal cleanup. This signal should not be used by users. It is also unwise to send `SIGKILL` to a server process — the main `postgres` process will interpret this as a crash and will force all the sibling processes to quit as part of its standard crash-recovery procedure.
Bugs

The -- options will not work on FreeBSD or OpenBSD. Use -c instead. This is a bug in the affected operating systems; a future release of PostgreSQL will provide a workaround if this is not fixed.

Single-User Mode

To start a single-user mode server, use a command like

```
postgres --single -D /usr/local/pgsql/data other-options my_database
```

Provide the correct path to the database directory with -D, or make sure that the environment variable PGDATA is set. Also specify the name of the particular database you want to work in.

Normally, the single-user mode server treats newline as the command entry terminator; there is no intelligence about semicolons, as there is in psql. To continue a command across multiple lines, you must type backslash just before each newline except the last one. The backslash and adjacent newline are both dropped from the input command. Note that this will happen even when within a string literal or comment.

But if you use the -j command line switch, a single newline does not terminate command entry; instead, the sequence semicolon-newline-newline does. That is, type a semicolon immediately followed by a completely empty line. Backslash-newline is not treated specially in this mode. Again, there is no intelligence about such a sequence appearing within a string literal or comment.

In either input mode, if you type a semicolon that is not just before or part of a command entry terminator, it is considered a command separator. When you do type a command entry terminator, the multiple statements you’ve entered will be executed as a single transaction.

To quit the session, type EOF (Control+D, usually). If you’ve entered any text since the last command entry terminator, then EOF will be taken as a command entry terminator, and another EOF will be needed to exit.

Note that the single-user mode server does not provide sophisticated line-editing features (no command history, for example). Single-user mode also does not do any background processing, such as automatic checkpoints or replication.

Examples

To start **postgres** in the background using default values, type:

```
$ nohup postgres >logfile 2>&1 < /dev/null &
```

To start **postgres** with a specific port, e.g. 1234:

```
$ postgres -p 1234
```

To connect to this server using **psql**, specify this port with the -p option:

```
$ psql -p 1234
```

or set the environment variable **PGPORT**:

```
$ export PGPORT=1234
```
Named run-time parameters can be set in either of these styles:

```bash
postgres -c work_mem=1234
postgres --work_mem=1234
```

Either form overrides whatever setting might exist for `work_mem` in `postgresql.conf`. Notice that underscores in parameter names can be written as either underscore or dash on the command line. Except for short-term experiments, it’s probably better practice to edit the setting in `postgresql.conf` than to rely on a command-line switch to set a parameter.

**See Also**

*initdb*, *pg_ctl*
postmaster

Name
postmaster — PostgreSQL database server

Synopsis
postmaster [option...]

Description
postmaster is a deprecated alias of postgres.

See Also
postgres
VII. Internals

This part contains assorted information that might be of use to PostgreSQL developers.
Chapter 49. Overview of PostgreSQL Internals

Author: This chapter originated as part of *Enhancement of the ANSI SQL Implementation of PostgreSQL*, Stefan Simkovics’ Master’s Thesis prepared at Vienna University of Technology under the direction of O.Univ.Prof.Dr. Georg Gottlob and Univ.Ass. Mag. Katrin Seyr.

This chapter gives an overview of the internal structure of the backend of PostgreSQL. After having read the following sections you should have an idea of how a query is processed. This chapter does not aim to provide a detailed description of the internal operation of PostgreSQL, as such a document would be very extensive. Rather, this chapter is intended to help the reader understand the general sequence of operations that occur within the backend from the point at which a query is received, to the point at which the results are returned to the client.

49.1. The Path of a Query

Here we give a short overview of the stages a query has to pass in order to obtain a result.

1. A connection from an application program to the PostgreSQL server has to be established. The application program transmits a query to the server and waits to receive the results sent back by the server.
2. The parser stage checks the query transmitted by the application program for correct syntax and creates a query tree.
3. The rewrite system takes the query tree created by the parser stage and looks for any rules (stored in the system catalogs) to apply to the query tree. It performs the transformations given in the rule bodies.

One application of the rewrite system is in the realization of views. Whenever a query against a view (i.e., a virtual table) is made, the rewrite system rewrites the user’s query to a query that accesses the base tables given in the view definition instead.
4. The planner/optimizer takes the (rewritten) query tree and creates a query plan that will be the input to the executor.

It does so by first creating all possible paths leading to the same result. For example if there is an index on a relation to be scanned, there are two paths for the scan. One possibility is a simple sequential scan and the other possibility is to use the index. Next the cost for the execution of each path is estimated and the cheapest path is chosen. The cheapest path is expanded into a complete plan that the executor can use.
5. The executor recursively steps through the plan tree and retrieves rows in the way represented by the plan. The executor makes use of the storage system while scanning relations, performs sorts and joins, evaluates qualifications and finally hands back the rows derived.

In the following sections we will cover each of the above listed items in more detail to give a better understanding of PostgreSQL’s internal control and data structures.
**49.2. How Connections are Established**

PostgreSQL is implemented using a simple “process per user” client/server model. In this model there is one client process connected to exactly one server process. As we do not know ahead of time how many connections will be made, we have to use a master process that spawns a new server process every time a connection is requested. This master process is called postgres and listens at a specified TCP/IP port for incoming connections. Whenever a request for a connection is detected the postgres process spawns a new server process. The server tasks communicate with each other using semaphores and shared memory to ensure data integrity throughout concurrent data access.

The client process can be any program that understands the PostgreSQL protocol described in Chapter 51. Many clients are based on the C-language library libpq, but several independent implementations of the protocol exist, such as the Java JDBC driver.

Once a connection is established the client process can send a query to the backend (server). The query is transmitted using plain text, i.e., there is no parsing done in the frontend (client). The server parses the query, creates an execution plan, executes the plan and returns the retrieved rows to the client by transmitting them over the established connection.

**49.3. The Parser Stage**

The parser stage consists of two parts:

- The parser defined in gram.y and scan.l is built using the Unix tools bison and flex.
- The transformation process does modifications and augmentations to the data structures returned by the parser.

**49.3.1. Parser**

The parser has to check the query string (which arrives as plain text) for valid syntax. If the syntax is correct a parse tree is built up and handed back; otherwise an error is returned. The parser and lexer are implemented using the well-known Unix tools bison and flex.

The lexer is defined in the file scan.l and is responsible for recognizing identifiers, the SQL key words etc. For every key word or identifier that is found, a token is generated and handed to the parser.

The parser is defined in the file gram.y and consists of a set of grammar rules and actions that are executed whenever a rule is fired. The code of the actions (which is actually C code) is used to build up the parse tree.

The file scan.l is transformed to the C source file scan.c using the program flex and gram.y is transformed to gram.c using bison. After these transformations have taken place a normal C compiler can be used to create the parser. Never make any changes to the generated C files as they will be overwritten the next time flex or bison is called.

**Note:** The mentioned transformations and compilations are normally done automatically using the makefiles shipped with the PostgreSQL source distribution.
A detailed description of bison or the grammar rules given in gram.y would be beyond the scope of this paper. There are many books and documents dealing with flex and bison. You should be familiar with bison before you start to study the grammar given in gram.y otherwise you won’t understand what happens there.

### 49.3.2. Transformation Process

The parser stage creates a parse tree using only fixed rules about the syntactic structure of SQL. It does not make any lookups in the system catalogs, so there is no possibility to understand the detailed semantics of the requested operations. After the parser completes, the transformation process takes the tree handed back by the parser as input and does the semantic interpretation needed to understand which tables, functions, and operators are referenced by the query. The data structure that is built to represent this information is called the query tree.

The reason for separating raw parsing from semantic analysis is that system catalog lookups can only be done within a transaction, and we do not wish to start a transaction immediately upon receiving a query string. The raw parsing stage is sufficient to identify the transaction control commands (BEGIN, ROLLBACK, etc), and these can then be correctly executed without any further analysis. Once we know that we are dealing with an actual query (such as SELECT or UPDATE), it is okay to start a transaction if we’re not already in one. Only then can the transformation process be invoked.

The query tree created by the transformation process is structurally similar to the raw parse tree in most places, but it has many differences in detail. For example, a FuncCall node in the parse tree represents something that looks syntactically like a function call. This might be transformed to either a FuncExpr or Aggref node depending on whether the referenced name turns out to be an ordinary function or an aggregate function. Also, information about the actual data types of columns and expression results is added to the query tree.

### 49.4. The PostgreSQL Rule System

PostgreSQL supports a powerful rule system for the specification of views and ambiguous view updates. Originally the PostgreSQL rule system consisted of two implementations:

- The first one worked using row level processing and was implemented deep in the executor. The rule system was called whenever an individual row had been accessed. This implementation was removed in 1995 when the last official release of the Berkeley Postgres project was transformed into Postgres95.
- The second implementation of the rule system is a technique called query rewriting. The rewrite system is a module that exists between the parser stage and the planner/optimizer. This technique is still implemented.

The query rewiter is discussed in some detail in Chapter 39, so there is no need to cover it here. We will only point out that both the input and the output of the rewiter are query trees, that is, there is no change in the representation or level of semantic detail in the trees. Rewriting can be thought of as a form of macro expansion.
49.5. Planner/Optimizer

The task of the planner/optimizer is to create an optimal execution plan. A given SQL query (and hence, a query tree) can be actually executed in a wide variety of different ways, each of which will produce the same set of results. If it is computationally feasible, the query optimizer will examine each of these possible execution plans, ultimately selecting the execution plan that is expected to run the fastest.

Note: In some situations, examining each possible way in which a query can be executed would take an excessive amount of time and memory space. In particular, this occurs when executing queries involving large numbers of join operations. In order to determine a reasonable (not necessarily optimal) query plan in a reasonable amount of time, PostgreSQL uses a Genetic Query Optimizer (see Chapter 58) when the number of joins exceeds a threshold (see geqo_threshold).

The planner’s search procedure actually works with data structures called paths, which are simply cut-down representations of plans containing only as much information as the planner needs to make its decisions. After the cheapest path is determined, a full-fledged plan tree is built to pass to the executor. This represents the desired execution plan in sufficient detail for the executor to run it. In the rest of this section we’ll ignore the distinction between paths and plans.

49.5.1. Generating Possible Plans

The planner/optimizer starts by generating plans for scanning each individual relation (table) used in the query. The possible plans are determined by the available indexes on each relation. There is always the possibility of performing a sequential scan on a relation, so a sequential scan plan is always created. Assume an index is defined on a relation (for example a B-tree index) and a query contains the restriction relation.attribute OPR constant. If relation.attribute happens to match the key of the B-tree index and OPR is one of the operators listed in the index’s operator class, another plan is created using the B-tree index to scan the relation. If there are further indexes present and the restrictions in the query happen to match a key of an index, further plans will be considered. Index scan plans are also generated for indexes that have a sort ordering that can match the query’s ORDER BY clause (if any), or a sort ordering that might be useful for merge joining (see below).

If the query requires joining two or more relations, plans for joining relations are considered after all feasible plans have been found for scanning single relations. The three available join strategies are:

- **nested loop join**: The right relation is scanned once for every row found in the left relation. This strategy is easy to implement but can be very time consuming. (However, if the right relation can be scanned with an index scan, this can be a good strategy. It is possible to use values from the current row of the left relation as keys for the index scan of the right.)

- **merge join**: Each relation is sorted on the join attributes before the join starts. Then the two relations are scanned in parallel, and matching rows are combined to form join rows. This kind of join is more attractive because each relation has to be scanned only once. The required sorting might be achieved either by an explicit sort step, or by scanning the relation in the proper order using an index on the join key.

- **hash join**: the right relation is first scanned and loaded into a hash table, using its join attributes as hash keys. Next the left relation is scanned and the appropriate values of every row found are used as hash keys to locate the matching rows in the table.
Chapter 49. Overview of PostgreSQL Internals

When the query involves more than two relations, the final result must be built up by a tree of join steps, each with two inputs. The planner examines different possible join sequences to find the cheapest one.

If the query uses fewer than \texttt{geqo\_threshold} relations, a near-exhaustive search is conducted to find the best join sequence. The planner preferentially considers joins between any two relations for which there exist a corresponding join clause in the \textbf{WHERE} qualification (i.e., for which a restriction like \texttt{where \texttt{rel1\_attr1=rel2\_attr2}} exists). Join pairs with no join clause are considered only when there is no other choice, that is, a particular relation has no available join clauses to any other relation. All possible plans are generated for every join pair considered by the planner, and the one that is (estimated to be) the cheapest is chosen.

When \texttt{geqo\_threshold} is exceeded, the join sequences considered are determined by heuristics, as described in Chapter 58. Otherwise the process is the same.

The finished plan tree consists of sequential or index scans of the base relations, plus nested-loop, merge, or hash join nodes as needed, plus any auxiliary steps needed, such as sort nodes or aggregate-function calculation nodes. Most of these plan node types have the additional ability to do selection (discarding rows that do not meet a specified Boolean condition) and projection (computation of a derived column set based on given column values, that is, evaluation of scalar expressions where needed). One of the responsibilities of the planner is to attach selection conditions from the \textbf{WHERE} clause and computation of required output expressions to the most appropriate nodes of the plan tree.

49.6. Executor

The \texttt{executor} takes the plan created by the planner/optimizer and recursively processes it to extract the required set of rows. This is essentially a demand-pull pipeline mechanism. Each time a plan node is called, it must deliver one more row, or report that it is done delivering rows.

To provide a concrete example, assume that the top node is a \texttt{MergeJoin} node. Before any merge can be done two rows have to be fetched (one from each subplan). So the executor recursively calls itself to process the subplans (it starts with the subplan attached to \texttt{left\_tree}). The new top node (the top node of the left subplan) is, let’s say, a \texttt{Sort} node and again recursion is needed to obtain an input row. The child node of the \texttt{Sort} might be a \texttt{SeqScan} node, representing actual reading of a table. Execution of this node causes the executor to fetch a row from the table and return it up to the calling node. The \texttt{Sort} node will repeatedly call its child to obtain all the rows to be sorted. When the input is exhausted (as indicated by the child node returning a NULL instead of a row), the \texttt{Sort} code performs the sort, and finally is able to return its first output row, namely the first one in sorted order. It keeps the remaining rows stored so that it can deliver them in sorted order in response to later demands.

The \texttt{MergeJoin} node similarly demands the first row from its right subplan. Then it compares the two rows to see if they can be joined; if so, it returns a join row to its caller. On the next call, or immediately if it cannot join the current pair of inputs, it advances to the next row of one table or the other (depending on how the comparison came out), and again checks for a match. Eventually, one subplan or the other is exhausted, and the \texttt{MergeJoin} node returns NULL to indicate that no more join rows can be formed.

Complex queries can involve many levels of plan nodes, but the general approach is the same: each node computes and returns its next output row each time it is called. Each node is also responsible for applying any selection or projection expressions that were assigned to it by the planner.

The executor mechanism is used to evaluate all four basic SQL query types: \texttt{SELECT}, \texttt{INSERT}, \texttt{UPDATE}, and \texttt{DELETE}. For \texttt{SELECT}, the top-level executor code only needs to send each row returned
by the query plan tree off to the client. For INSERT, each returned row is inserted into the target table specified for the INSERT. This is done in a special top-level plan node called ModifyTable. (A simple INSERT ... VALUES command creates a trivial plan tree consisting of a single Result node, which computes just one result row, and ModifyTable above it to perform the insertion. But INSERT ...

SELECT can demand the full power of the executor mechanism.) For UPDATE, the planner arranges that each computed row includes all the updated column values, plus the TID (tuple ID, or row ID) of the original target row; this data is fed into a ModifyTable node, which uses the information to create a new updated row and mark the old row deleted. For DELETE, the only column that is actually returned by the plan is the TID, and the ModifyTable node simply uses the TID to visit each target row and mark it deleted.
Chapter 50. System Catalogs

The system catalogs are the place where a relational database management system stores schema metadata, such as information about tables and columns, and internal bookkeeping information. PostgreSQL’s system catalogs are regular tables. You can drop and recreate the tables, add columns, insert and update values, and severely mess up your system that way. Normally, one should not change the system catalogs by hand, there are normally SQL commands to do that. (For example, CREATE DATABASE inserts a row into the pg_database catalog — and actually creates the database on disk.) There are some exceptions for particularly esoteric operations, but many of those have been made available as SQL commands over time, and so the need for direct manipulation of the system catalogs is ever decreasing.

50.1. Overview

Table 50-1 lists the system catalogs. More detailed documentation of each catalog follows below.

Most system catalogs are copied from the template database during database creation and are thereafter database-specific. A few catalogs are physically shared across all databases in a cluster; these are noted in the descriptions of the individual catalogs.

Table 50-1. System Catalogs

<table>
<thead>
<tr>
<th>Catalog Name</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_aggregate</td>
<td>aggregate functions</td>
</tr>
<tr>
<td>pg_am</td>
<td>index access methods</td>
</tr>
<tr>
<td>pg_amop</td>
<td>access method operators</td>
</tr>
<tr>
<td>pg_amproc</td>
<td>access method support procedures</td>
</tr>
<tr>
<td>pg_attrdef</td>
<td>column default values</td>
</tr>
<tr>
<td>pg_attribute</td>
<td>table columns (“attributes”)</td>
</tr>
<tr>
<td>pg_authid</td>
<td>authorization identifiers (roles)</td>
</tr>
<tr>
<td>pg_auth_members</td>
<td>authorization identifier membership relationships</td>
</tr>
<tr>
<td>pg_cast</td>
<td>casts (data type conversions)</td>
</tr>
<tr>
<td>pg_class</td>
<td>tables, indexes, sequences, views (“relations”)</td>
</tr>
<tr>
<td>pg_collation</td>
<td>collations (locale information)</td>
</tr>
<tr>
<td>pg_constraint</td>
<td>check constraints, unique constraints, primary key constraints, foreign key constraints</td>
</tr>
<tr>
<td>pg_conversion</td>
<td>encoding conversion information</td>
</tr>
<tr>
<td>pg_database</td>
<td>databases within this database cluster</td>
</tr>
<tr>
<td>pg_db_role_setting</td>
<td>per-role and per-database settings</td>
</tr>
<tr>
<td>pg_default_acl</td>
<td>default privileges for object types</td>
</tr>
<tr>
<td>pg_depend</td>
<td>dependencies between database objects</td>
</tr>
<tr>
<td>pg_description</td>
<td>descriptions or comments on database objects</td>
</tr>
<tr>
<td>Catalog Name</td>
<td>Purpose</td>
</tr>
<tr>
<td>--------------------------</td>
<td>----------------------------------------------</td>
</tr>
<tr>
<td>pg_enum</td>
<td>enum label and value definitions</td>
</tr>
<tr>
<td>pg_event_trigger</td>
<td>event triggers</td>
</tr>
<tr>
<td>pg_extension</td>
<td>installed extensions</td>
</tr>
<tr>
<td>pg_foreign_data_wrapper</td>
<td>foreign-data wrapper definitions</td>
</tr>
<tr>
<td>pg_foreign_server</td>
<td>foreign server definitions</td>
</tr>
<tr>
<td>pg_foreign_table</td>
<td>additional foreign table information</td>
</tr>
<tr>
<td>pg_index</td>
<td>additional index information</td>
</tr>
<tr>
<td>pg_inherits</td>
<td>table inheritance hierarchy</td>
</tr>
<tr>
<td>pg_init_privs</td>
<td>object initial privileges</td>
</tr>
<tr>
<td>pg_language</td>
<td>languages for writing functions</td>
</tr>
<tr>
<td>pg_largeobject</td>
<td>data pages for large objects</td>
</tr>
<tr>
<td>pg_largeobject_metadata</td>
<td>metadata for large objects</td>
</tr>
<tr>
<td>pg_namespace</td>
<td>schemas</td>
</tr>
<tr>
<td>pg_opclass</td>
<td>access method operator classes</td>
</tr>
<tr>
<td>pg_operator</td>
<td>operators</td>
</tr>
<tr>
<td>pg_opfamily</td>
<td>access method operator families</td>
</tr>
<tr>
<td>pg_pltemplate</td>
<td>template data for procedural languages</td>
</tr>
<tr>
<td>pg_policy</td>
<td>row-security policies</td>
</tr>
<tr>
<td>pg_proc</td>
<td>functions and procedures</td>
</tr>
<tr>
<td>pg_range</td>
<td>information about range types</td>
</tr>
<tr>
<td>pg_replication_origin</td>
<td>registered replication origins</td>
</tr>
<tr>
<td>pg_rewrite</td>
<td>query rewrite rules</td>
</tr>
<tr>
<td>pg_seclabel</td>
<td>security labels on database objects</td>
</tr>
<tr>
<td>pg_shdepend</td>
<td>dependencies on shared objects</td>
</tr>
<tr>
<td>pg_shdescription</td>
<td>comments on shared objects</td>
</tr>
<tr>
<td>pg_shseclabel</td>
<td>security labels on shared database objects</td>
</tr>
<tr>
<td>pg_statistic</td>
<td>planner statistics</td>
</tr>
<tr>
<td>pg_tablespace</td>
<td>table spaces within this database cluster</td>
</tr>
<tr>
<td>pg_transform</td>
<td>transforms (data type to procedural language conversions)</td>
</tr>
<tr>
<td>pg_trigger</td>
<td>triggers</td>
</tr>
<tr>
<td>pg_ts_config</td>
<td>text search configurations</td>
</tr>
<tr>
<td>pg_ts_config_map</td>
<td>text search configurations’ token mappings</td>
</tr>
<tr>
<td>pg_ts_dict</td>
<td>text search dictionaries</td>
</tr>
<tr>
<td>pg_ts_parser</td>
<td>text search parsers</td>
</tr>
<tr>
<td>pg_ts_template</td>
<td>text search templates</td>
</tr>
<tr>
<td>pg_type</td>
<td>data types</td>
</tr>
<tr>
<td>pg_user_mapping</td>
<td>mappings of users to foreign servers</td>
</tr>
</tbody>
</table>
50.2. pg_aggregate

The catalog pg_aggregate stores information about aggregate functions. An aggregate function is a function that operates on a set of values (typically one column from each row that matches a query condition) and returns a single value computed from all these values. Typical aggregate functions are `sum`, `count`, and `max`. Each entry in pg_aggregate is an extension of an entry in pg_proc. The pg_proc entry carries the aggregate’s name, input and output data types, and other information that is similar to ordinary functions.

Table 50-2. pg_aggregate Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>aggfnoid</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>pg_proc OID of the aggregate function</td>
</tr>
<tr>
<td>aggkind</td>
<td>char</td>
<td></td>
<td>Aggregate kind: n for “normal” aggregates, o for “ordered-set” aggregates, or h for “hypothetical-set” aggregates</td>
</tr>
<tr>
<td>aggnumdirectargs</td>
<td>int2</td>
<td></td>
<td>Number of direct (non-aggregated) arguments of an ordered-set or hypothetical-set aggregate, counting a variadic array as one argument. If equal to pronargs, the aggregate must be variadic and the variadic array describes the aggregated arguments as well as the final direct arguments. Always zero for normal aggregates.</td>
</tr>
<tr>
<td>aggtransfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Transition function</td>
</tr>
<tr>
<td>aggfinalfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Final function (zero if none)</td>
</tr>
<tr>
<td>aggcombinefn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Combine function (zero if none)</td>
</tr>
<tr>
<td>aggserialfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Serialization function (zero if none)</td>
</tr>
<tr>
<td>aggsdeserialfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Deserialization function (zero if none)</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>--------</td>
<td>----------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>aggmtransfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Forward transition function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggminvtransfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Inverse transition function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggmfinalfn</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Final function for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggfinalextra</td>
<td>bool</td>
<td></td>
<td>True to pass extra dummy arguments to aggmfinalfn</td>
</tr>
<tr>
<td>aggmfinalfull</td>
<td>bool</td>
<td></td>
<td>True to pass extra dummy arguments to aggmfinalfn</td>
</tr>
<tr>
<td>aggsortop</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>Associated sort operator (zero if none)</td>
</tr>
<tr>
<td>aggtranstype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type of the aggregate function’s internal transition (state) data</td>
</tr>
<tr>
<td>aggmtransspace</td>
<td>int4</td>
<td></td>
<td>Approximate average size (in bytes) of the transition state data, or zero to use a default estimate</td>
</tr>
<tr>
<td>aggmtransspace</td>
<td>int4</td>
<td></td>
<td>Data type of the aggregate function’s internal transition (state) data for moving-aggregate mode (zero if none)</td>
</tr>
<tr>
<td>aggmtransspace</td>
<td>int4</td>
<td></td>
<td>Approximate average size (in bytes) of the transition state data for moving-aggregate mode, or zero to use a default estimate</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

#### 50.3. `pg_am`

The catalog `pg_am` stores information about relation access methods. There is one row for each access method supported by the system. Currently, only indexes have access methods. The requirements for index access methods are discussed in detail in Chapter 59.

#### Table 50-3. `pg_am` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>amname</td>
<td>name</td>
<td></td>
<td>Name of the access method</td>
</tr>
<tr>
<td>amhandler</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of a handler function that is responsible for supplying information about the access method</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>amtype</td>
<td>char</td>
<td>i</td>
<td>Currently always i to indicate an index access method; other values may be allowed in future</td>
</tr>
</tbody>
</table>

**Note:** Before PostgreSQL 9.6, pg_am contained many additional columns representing properties of index access methods. That data is now only directly visible at the C code level. However, pg_index_column_has_property() and related functions have been added to allow SQL queries to inspect index access method properties; see Table 9-62.

#### 50.4. pg_amop

The catalog pg_amop stores information about operators associated with access method operator families. There is one row for each operator that is a member of an operator family. A family member can be either a search operator or an ordering operator. An operator can appear in more than one family, but cannot appear in more than one search position nor more than one ordering position within a family. (It is allowed, though unlikely, for an operator to be used for both search and ordering purposes.)

**Table 50-4. pg_amop Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>amopfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>The operator family this entry is for</td>
</tr>
<tr>
<td>amoplefttype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Left-hand input data type of operator</td>
</tr>
<tr>
<td>amoprighttype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Right-hand input data type of operator</td>
</tr>
<tr>
<td>amopstrategy</td>
<td>int2</td>
<td></td>
<td>Operator strategy number</td>
</tr>
<tr>
<td>amoppurpose</td>
<td>char</td>
<td></td>
<td>Operator purpose, either s for search or o for ordering</td>
</tr>
<tr>
<td>amopopr</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>OID of the operator</td>
</tr>
<tr>
<td>amopmethod</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>Index access method operator family is for</td>
</tr>
<tr>
<td>amopsortfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>The B-tree operator family this entry sorts according to, if an ordering operator; zero if a search operator</td>
</tr>
</tbody>
</table>
A “search” operator entry indicates that an index of this operator family can be searched to find all rows satisfying \( \text{WHERE indexed_column operator constant} \). Obviously, such an operator must return boolean, and its left-hand input type must match the index’s column data type.

An “ordering” operator entry indicates that an index of this operator family can be scanned to return rows in the order represented by \( \text{ORDER BY indexed_column operator constant} \). Such an operator could return any sortable data type, though again its left-hand input type must match the index’s column data type. The exact semantics of the \( \text{ORDER BY} \) are specified by the amopsortfamily column, which must reference a B-tree operator family for the operator’s result type.

**Note:** At present, it’s assumed that the sort order for an ordering operator is the default for the referenced operator family, i.e., ASC NULLS LAST. This might someday be relaxed by adding additional columns to specify sort options explicitly.

An entry’s amopmethod must match the opmethod of its containing operator family (including amopmethod here is an intentional denormalization of the catalog structure for performance reasons). Also, amoplefttype and amoprighttype must match the oprleft and oprright fields of the referenced pg_operator entry.

### 50.5. pg_amproc

The catalog pg_amproc stores information about support procedures associated with access method operator families. There is one row for each support procedure belonging to an operator family.

#### Table 50-5. pg_amproc Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>amprocfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>The operator family this entry is for</td>
</tr>
<tr>
<td>amproclefttype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Left-hand input data type of associated operator</td>
</tr>
<tr>
<td>amprocrighttype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Right-hand input data type of associated operator</td>
</tr>
<tr>
<td>amprocnum</td>
<td>int2</td>
<td></td>
<td>Support procedure number</td>
</tr>
<tr>
<td>amproc</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the procedure</td>
</tr>
</tbody>
</table>

The usual interpretation of the amproclefttype and amprocrighttype fields is that they identify the left and right input types of the operator(s) that a particular support procedure supports. For some access methods these match the input data type(s) of the support procedure itself, for others not. There is a notion of “default” support procedures for an index, which are those with amproclefttype and amprocrighttype both equal to the index operator class’s opcintype.
50.6. `pg_attrdef`

The catalog `pg_attrdef` stores column default values. The main information about columns is stored in `pg_attribute` (see below). Only columns that explicitly specify a default value (when the table is created or the column is added) will have an entry here.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>adrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this column belongs to</td>
</tr>
<tr>
<td>adnum</td>
<td>int2</td>
<td>pg_attribute.attnum</td>
<td>The number of the column</td>
</tr>
<tr>
<td>adbin</td>
<td>pg_node_tree</td>
<td></td>
<td>The internal representation of the column default value</td>
</tr>
<tr>
<td>adsrc</td>
<td>text</td>
<td></td>
<td>A human-readable representation of the default value</td>
</tr>
</tbody>
</table>

The `adsrc` field is historical, and is best not used, because it does not track outside changes that might affect the representation of the default value. Reverse-compiling the `adbin` field (with `pg_get_expr` for example) is a better way to display the default value.

50.7. `pg_attribute`

The catalog `pg_attribute` stores information about table columns. There will be exactly one `pg_attribute` row for every column in every table in the database. (There will also be attribute entries for indexes, and indeed all objects that have `pg_class` entries.)

The term attribute is equivalent to column and is used for historical reasons.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>attrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this column belongs to</td>
</tr>
<tr>
<td>attname</td>
<td>name</td>
<td></td>
<td>The column name</td>
</tr>
<tr>
<td>atttypid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>The data type of this column</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>---------------</td>
<td>------</td>
<td>------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>attstattarget</td>
<td>int4</td>
<td></td>
<td>attstattarget controls the level of detail of statistics accumulated for this column by ANALYZE. A zero value indicates that no statistics should be collected. A negative value says to use the system default statistics target. The exact meaning of positive values is data type-dependent. For scalar data types, attstattarget is both the target number of “most common values” to collect, and the target number of histogram bins to create.</td>
</tr>
<tr>
<td>attlen</td>
<td>int2</td>
<td></td>
<td>A copy of pg_type.typlen of this column’s type</td>
</tr>
<tr>
<td>attnum</td>
<td>int2</td>
<td></td>
<td>The number of the column. Ordinary columns are numbered from 1 up. System columns, such as oid, have (arbitrary) negative numbers.</td>
</tr>
<tr>
<td>attndims</td>
<td>int4</td>
<td></td>
<td>Number of dimensions, if the column is an array type; otherwise 0. (Presently, the number of dimensions of an array is not enforced, so any nonzero value effectively means “it’s an array”.)</td>
</tr>
<tr>
<td>attcacheoff</td>
<td>int4</td>
<td></td>
<td>Always -1 in storage, but when loaded into a row descriptor in memory this might be updated to cache the offset of the attribute within the row</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>---------</td>
<td>------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>atttypmod</td>
<td>int4</td>
<td></td>
<td><code>atttypmod</code> records type-specific data supplied at table creation time (for example, the maximum length of a <code>varchar</code> column). It is passed to type-specific input functions and length coercion functions. The value will generally be -1 for types that do not need <code>atttypmod</code>.</td>
</tr>
<tr>
<td>attbyval</td>
<td>bool</td>
<td></td>
<td>A copy of <code>pg_type.typbyval</code> of this column’s type.</td>
</tr>
<tr>
<td>attstorage</td>
<td>char</td>
<td></td>
<td>Normally a copy of <code>pg_type.typstorage</code> of this column’s type. For TOAST-able data types, this can be altered after column creation to control storage policy.</td>
</tr>
<tr>
<td>attalign</td>
<td>char</td>
<td></td>
<td>A copy of <code>pg_type.typalign</code> of this column’s type.</td>
</tr>
<tr>
<td>attnotnull</td>
<td>bool</td>
<td></td>
<td>This represents a not-null constraint.</td>
</tr>
<tr>
<td>atthasdef</td>
<td>bool</td>
<td></td>
<td>This column has a default value, in which case there will be a corresponding entry in the <code>pg_attrdef</code> catalog that actually defines the value.</td>
</tr>
<tr>
<td>attisdropped</td>
<td>bool</td>
<td></td>
<td>This column has been dropped and is no longer valid. A dropped column is still physically present in the table, but is ignored by the parser and so cannot be accessed via SQL.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>----------------</td>
<td>-----------------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>attislocal</td>
<td>bool</td>
<td></td>
<td>This column is defined locally in the relation. Note that a column can be</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>locally defined and inherited simultaneously.</td>
</tr>
<tr>
<td>attinhcount</td>
<td>int4</td>
<td></td>
<td>The number of direct ancestors this column has. A column with a nonzero</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>number of ancestors cannot be dropped nor renamed.</td>
</tr>
<tr>
<td>attcollation</td>
<td>oid</td>
<td>pg_collation.oid</td>
<td>The defined collation of the column, or zero if the column is not of a</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>collatable data type.</td>
</tr>
<tr>
<td>attacl</td>
<td>aclitem[]</td>
<td></td>
<td>Column-level access privileges, if any have been granted specifically on this</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>column.</td>
</tr>
<tr>
<td>attoptions</td>
<td>text[]</td>
<td></td>
<td>Attribute-level options, as “keyword=value” strings</td>
</tr>
<tr>
<td>attfdwoptions</td>
<td>text[]</td>
<td></td>
<td>Attribute-level foreign data wrapper options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

In a dropped column’s `pg_attribute` entry, `atttypid` is reset to zero, but `attlen` and the other fields copied from `pg_type` are still valid. This arrangement is needed to cope with the situation where the dropped column’s data type was later dropped, and so there is no `pg_type` row anymore. `attlen` and the other fields can be used to interpret the contents of a row of the table.

### 50.8. `pg_authid`

The catalog `pg_authid` contains information about database authorization identifiers (roles). A role subsumes the concepts of “users” and “groups”. A user is essentially just a role with the `rolcanlogin` flag set. Any role (with or without `rolcanlogin`) can have other roles as members; see `pg_auth_members`.

Since this catalog contains passwords, it must not be publicly readable. `pg_roles` is a publicly readable view on `pg_authid` that blanks out the password field.

Chapter 21 contains detailed information about user and privilege management.

Because user identities are cluster-wide, `pg_authid` is shared across all databases of a cluster: there is only one copy of `pg_authid` per cluster, not one per database.
### Table 50-8. pg_authid Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>rolname</td>
<td>name</td>
<td>Role name</td>
</tr>
<tr>
<td>rolsuper</td>
<td>bool</td>
<td>Role has superuser privileges</td>
</tr>
<tr>
<td>rolinherit</td>
<td>bool</td>
<td>Role automatically inherits privileges of roles it is a member of</td>
</tr>
<tr>
<td>rolcreaterole</td>
<td>bool</td>
<td>Role can create more roles</td>
</tr>
<tr>
<td>rolcreatedb</td>
<td>bool</td>
<td>Role can create databases</td>
</tr>
<tr>
<td>rolcanlogin</td>
<td>bool</td>
<td>Role can log in. That is, this role can be given as the initial session authorization identifier</td>
</tr>
<tr>
<td>rolreplication</td>
<td>bool</td>
<td>Role is a replication role. A replication role can initiate replication connections and create and drop replication slots.</td>
</tr>
<tr>
<td>rolbypassrls</td>
<td>bool</td>
<td>Role bypasses every row level security policy, see Section 5.7 for more information.</td>
</tr>
<tr>
<td>rolconnlimit</td>
<td>int4</td>
<td>For roles that can log in, this sets maximum number of concurrent connections this role can make. -1 means no limit.</td>
</tr>
<tr>
<td>rolpassword</td>
<td>text</td>
<td>Password (possibly encrypted); null if none. If the password is encrypted, this column will begin with the string \texttt{md5} followed by a 32-character hexadecimal MD5 hash. The MD5 hash will be of the user’s password concatenated to their user name. For example, if user joe has password xyzzy, PostgreSQL will store the md5 hash of xyzzyjoe. A password that does not follow that format is assumed to be unencrypted.</td>
</tr>
<tr>
<td>rolvaliduntil</td>
<td>timestamptz</td>
<td>Password expiry time (only used for password authentication); null if no expiration</td>
</tr>
</tbody>
</table>
50.9. pg_auth_members

The catalog `pg_auth_members` shows the membership relations between roles. Any non-circular set of relationships is allowed.

Because user identities are cluster-wide, `pg_auth_members` is shared across all databases of a cluster: there is only one copy of `pg_auth_members` per cluster, not one per database.

Table 50-9. pg_auth_members Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>roleid</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of a role that has a member</td>
</tr>
<tr>
<td>member</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of a role that is a member of roleid</td>
</tr>
<tr>
<td>grantor</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of the role that granted this membership</td>
</tr>
<tr>
<td>admin_option</td>
<td>bool</td>
<td></td>
<td>True if member can grant membership in roleid to others</td>
</tr>
</tbody>
</table>

50.10. pg_cast

The catalog `pg_cast` stores data type conversion paths, both built-in and user-defined.

It should be noted that `pg_cast` does not represent every type conversion that the system knows how to perform; only those that cannot be deduced from some generic rule. For example, casting between a domain and its base type is not explicitly represented in `pg_cast`. Another important exception is that “automatic I/O conversion casts”, those performed using a data type’s own I/O functions to convert to or from `text` or other string types, are not explicitly represented in `pg_cast`.

Table 50-10. pg_cast Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>castsoure</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the source data type</td>
</tr>
<tr>
<td>casttarget</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the target data type</td>
</tr>
<tr>
<td>castfunc</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>The OID of the function to use to perform this cast. Zero is stored if the cast method doesn’t require a function.</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

### Table 50-11. pg_cast Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>castcontext</td>
<td>char</td>
<td></td>
<td>Indicates what contexts the cast can be invoked in. e means only as an explicit cast (using CAST or :: syntax), a means implicitly in assignment to a target column, as well as explicitly. i means implicitly in expressions, as well as the other cases.</td>
</tr>
<tr>
<td>castmethod</td>
<td>char</td>
<td></td>
<td>Indicates how the cast is performed. f means that the function specified in the castfunc field is used. i means that the input/output functions are used. b means that the types are binary-coercible, thus no conversion is required.</td>
</tr>
</tbody>
</table>

The cast functions listed in pg_cast must always take the cast source type as their first argument type, and return the cast destination type as their result type. A cast function can have up to three arguments. The second argument, if present, must be type integer; it receives the type modifier associated with the destination type, or -1 if there is none. The third argument, if present, must be type boolean; it receives true if the cast is an explicit cast, false otherwise.

It is legitimate to create a pg_cast entry in which the source and target types are the same, if the associated function takes more than one argument. Such entries represent “length coercion functions” that coerce values of the type to be legal for a particular type modifier value.

When a pg_cast entry has different source and target types and a function that takes more than one argument, it represents converting from one type to another and applying a length coercion in a single step. When no such entry is available, coercion to a type that uses a type modifier involves two steps, one to convert between data types and a second to apply the modifier.

### 50.11. pg_class

The catalog pg_class catalogs tables and most everything else that has columns or is otherwise similar to a table. This includes indexes (but see also pg_index), sequences, views, materialized views, composite types, and TOAST tables; see relkind. Below, when we mean all of these kinds of objects we speak of “relations”. Not all columns are meaningful for all relation types.

<p>| Table 50-11. pg_class Columns |</p>
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>relname</td>
<td>name</td>
<td></td>
<td>Name of the table, index, view, etc.</td>
</tr>
<tr>
<td>relnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this relation</td>
</tr>
<tr>
<td>reltype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>The OID of the data type that corresponds to this table’s row type, if any (zero for indexes, which have no pg_type entry)</td>
</tr>
<tr>
<td>relloftype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>For typed tables, the OID of the underlying composite type, zero for all other relations</td>
</tr>
<tr>
<td>reloowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the relation</td>
</tr>
<tr>
<td>relam</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>If this is an index, the access method used (B-tree, hash, etc.)</td>
</tr>
<tr>
<td>relfilenode</td>
<td>oid</td>
<td></td>
<td>Name of the on-disk file of this relation; zero means this is a “mapped” relation whose disk file name is determined by low-level state</td>
</tr>
<tr>
<td>reltablespace</td>
<td>oid</td>
<td>pg_tablespace.oid</td>
<td>The tablespace in which this relation is stored. If zero, the database’s default tablespace is implied. (Not meaningful if the relation has no on-disk file.)</td>
</tr>
<tr>
<td>relpages</td>
<td>int4</td>
<td></td>
<td>Size of the on-disk representation of this table in pages (of size BLCKSZ). This is only an estimate used by the planner. It is updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE INDEX.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>------------</td>
<td>-------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>reltuples</td>
<td>float4</td>
<td></td>
<td>Number of rows in the table. This is only an estimate used by the planner. It is updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE_INDEX.</td>
</tr>
<tr>
<td>relallvisible</td>
<td>int4</td>
<td></td>
<td>Number of pages that are marked all-visible in the table’s visibility map. This is only an estimate used by the planner. It is updated by VACUUM, ANALYZE, and a few DDL commands such as CREATE_INDEX.</td>
</tr>
<tr>
<td>reltoastrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>OID of the TOAST table associated with this table. 0 if none. The TOAST table stores large attributes “out of line” in a secondary table.</td>
</tr>
<tr>
<td>relhasindex</td>
<td>bool</td>
<td></td>
<td>True if this is a table and it has (or recently had) any indexes</td>
</tr>
<tr>
<td>relisshared</td>
<td>bool</td>
<td></td>
<td>True if this table is shared across all databases in the cluster. Only certain system catalogs (such as pg_database) are shared.</td>
</tr>
<tr>
<td>relpersistence</td>
<td>char</td>
<td></td>
<td>p = permanent table, u = unlogged table, t = temporary table</td>
</tr>
<tr>
<td>relkind</td>
<td>char</td>
<td></td>
<td>r = ordinary table, i = index, s = sequence, v = view, m = materialized view, c = composite type, t = TOAST table, f = foreign table</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>-------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>relnatts</td>
<td>int2</td>
<td></td>
<td>Number of user columns in the relation (system columns not counted). There must be this many corresponding entries in pg_attribute. See also pg_attribute.attnum.</td>
</tr>
<tr>
<td>relchecks</td>
<td>int2</td>
<td></td>
<td>Number of CHECK constraints on the table; see pg_constraint catalog</td>
</tr>
<tr>
<td>relhasoids</td>
<td>bool</td>
<td></td>
<td>True if we generate an OID for each row of the relation</td>
</tr>
<tr>
<td>relhaspkey</td>
<td>bool</td>
<td></td>
<td>True if the table has (or once had) a primary key</td>
</tr>
<tr>
<td>relhasrules</td>
<td>bool</td>
<td></td>
<td>True if table has (or once had) rules; see pg_rewrite catalog</td>
</tr>
<tr>
<td>relhastriggers</td>
<td>bool</td>
<td></td>
<td>True if table has (or once had) triggers; see pg_trigger catalog</td>
</tr>
<tr>
<td>relhassubclass</td>
<td>bool</td>
<td></td>
<td>True if table has (or once had) any inheritance children</td>
</tr>
<tr>
<td>relrowsecurity</td>
<td>bool</td>
<td></td>
<td>True if table has row level security enabled; see pg_policy catalog</td>
</tr>
<tr>
<td>relforcerowsecurity</td>
<td>bool</td>
<td></td>
<td>True if row level security (when enabled) will also apply to table owner; see pg_policy catalog</td>
</tr>
<tr>
<td>relispopulated</td>
<td>bool</td>
<td></td>
<td>True if relation is populated (this is true for all relations other than some materialized views)</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>-------------</td>
<td>------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>relreplident</td>
<td>char</td>
<td></td>
<td>Columns used to form &quot;replica identity&quot; for rows: (d = ) default (primary key, if any), (n = ) nothing, (f = ) all columns (i = ) index with indisreplident set, or default</td>
</tr>
<tr>
<td>relfrozenxid</td>
<td>xid</td>
<td></td>
<td>All transaction IDs before this one have been replaced with a permanent (&quot;frozen&quot;) transaction ID in this table. This is used to track whether the table needs to be vacuumed in order to prevent transaction ID wraparound or to allow pg_clog to be shrunk. Zero (InvalidTransactionId) if the relation is not a table.</td>
</tr>
<tr>
<td>relminmixinid</td>
<td>xid</td>
<td></td>
<td>All multixact IDs before this one have been replaced by a transaction ID in this table. This is used to track whether the table needs to be vacuumed in order to prevent multixact ID wraparound or to allow pg_multixact to be shrunk. Zero (InvalidMultiXactId) if the relation is not a table.</td>
</tr>
<tr>
<td>relacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
<tr>
<td>reloptions</td>
<td>text[]</td>
<td></td>
<td>Access-method-specific options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

Several of the Boolean flags in pg_class are maintained lazily: they are guaranteed to be true if that’s the correct state, but may not be reset to false immediately when the condition is no longer true. For example, relhasindex is set by CREATE INDEX, but it is never cleared by DROP INDEX.
Instead, VACUUM clears relhasindex if it finds the table has no indexes. This arrangement avoids race conditions and improves concurrency.

### 50.12. pg_collation

The catalog pg_collation describes the available collations, which are essentially mappings from an SQL name to operating system locale categories. See Section 23.2 for more information.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>collname</td>
<td>name</td>
<td></td>
<td>Collation name (unique per namespace and encoding)</td>
</tr>
<tr>
<td>collnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this collation</td>
</tr>
<tr>
<td>collowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the collation</td>
</tr>
<tr>
<td>collencoding</td>
<td>int4</td>
<td></td>
<td>Encoding in which the collation is applicable, or -1 if it works for any encoding</td>
</tr>
<tr>
<td>collcollate</td>
<td>name</td>
<td></td>
<td>LC_COLLATE for this collation object</td>
</tr>
<tr>
<td>collctype</td>
<td>name</td>
<td></td>
<td>LC_CTYPE for this collation object</td>
</tr>
</tbody>
</table>

Note that the unique key on this catalog is (collname, collencoding, collnamespace) not just (collname, collnamespace). PostgreSQL generally ignores all collations that do not have collencoding equal to either the current database’s encoding or -1, and creation of new entries with the same name as an entry with collencoding = -1 is forbidden. Therefore it is sufficient to use a qualified SQL name (schema.name) to identify a collation, even though this is not unique according to the catalog definition. The reason for defining the catalog this way is that initdb fills it in at cluster initialization time with entries for all locales available on the system, so it must be able to hold entries for all encodings that might ever be used in the cluster.

In the template0 database, it could be useful to create collations whose encoding does not match the database encoding, since they could match the encodings of databases later cloned from template0. This would currently have to be done manually.

### 50.13. pg_constraint

The catalog pg_constraint stores check, primary key, unique, foreign key, and exclusion constraints on tables. (Column constraints are not treated specially. Every column constraint is equivalent to some table constraint.) Not-null constraints are represented in the pg_attribute catalog, not
User-defined constraint triggers (created with \texttt{CREATE CONSTRAINT TRIGGER}) also give rise to an entry in this table.

Check constraints on domains are stored here, too.

Table 50-13. \texttt{pg\_constraint} Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{oid}</td>
<td>\texttt{oid}</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>\texttt{conname}</td>
<td>\texttt{name}</td>
<td></td>
<td>Constraint name (not necessarily unique!)</td>
</tr>
<tr>
<td>\texttt{connamespace}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_namespace.oid}</td>
<td>The OID of the namespace that contains this constraint</td>
</tr>
<tr>
<td>\texttt{contype}</td>
<td>\texttt{char}</td>
<td></td>
<td>\texttt{c} = check constraint, \texttt{f} = foreign key constraint, \texttt{p} = primary key constraint, \texttt{u} = unique constraint, \texttt{t} = constraint trigger, \texttt{x} = exclusion constraint</td>
</tr>
<tr>
<td>\texttt{condeferrable}</td>
<td>\texttt{bool}</td>
<td></td>
<td>Is the constraint deferrable?</td>
</tr>
<tr>
<td>\texttt{condeferred}</td>
<td>\texttt{bool}</td>
<td></td>
<td>Is the constraint deferred by default?</td>
</tr>
<tr>
<td>\texttt{convalidated}</td>
<td>\texttt{bool}</td>
<td></td>
<td>Has the constraint been validated? Currently, can only be false for foreign keys and CHECK constraints</td>
</tr>
<tr>
<td>\texttt{conrelid}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_class.oid}</td>
<td>The table this constraint is on; 0 if not a table constraint</td>
</tr>
<tr>
<td>\texttt{contypid}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_type.oid}</td>
<td>The domain this constraint is on; 0 if not a domain constraint</td>
</tr>
<tr>
<td>\texttt{conindid}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_class.oid}</td>
<td>The index supporting this constraint, if it’s a unique, primary key, foreign key, or exclusion constraint; else 0</td>
</tr>
<tr>
<td>\texttt{confrelid}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_class.oid}</td>
<td>If a foreign key, the referenced table; else 0</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>---------------</td>
<td>-----------</td>
<td>------------</td>
<td>------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>confupdtype</td>
<td>char</td>
<td></td>
<td>Foreign key update action code: (a = \text{no action}, r = \text{restrict}, c = \text{cascade}, n = \text{set null}, d = \text{set default})</td>
</tr>
<tr>
<td>confdeltype</td>
<td>char</td>
<td></td>
<td>Foreign key deletion action code: (a = \text{no action}, r = \text{restrict}, c = \text{cascade}, n = \text{set null}, d = \text{set default})</td>
</tr>
<tr>
<td>confmatchtype</td>
<td>char</td>
<td></td>
<td>Foreign key match type: (f = \text{full}, p = \text{partial}, s = \text{simple})</td>
</tr>
<tr>
<td>conislocal</td>
<td>bool</td>
<td></td>
<td>This constraint is defined locally for the relation. Note that a constraint can be locally defined and inherited simultaneously.</td>
</tr>
<tr>
<td>coninhcount</td>
<td>int4</td>
<td></td>
<td>The number of direct inheritance ancestors this constraint has. A constraint with a nonzero number of ancestors cannot be dropped nor renamed.</td>
</tr>
<tr>
<td>connoinherit</td>
<td>bool</td>
<td></td>
<td>This constraint is defined locally for the relation. It is a non-inheritable constraint.</td>
</tr>
<tr>
<td>conkey</td>
<td>int2[]</td>
<td>pg_attribute.attnum</td>
<td>If a table constraint (including foreign keys, but not constraint triggers), list of the constrained columns</td>
</tr>
<tr>
<td>confkey</td>
<td>int2[]</td>
<td>pg_attribute.attnum</td>
<td>If a foreign key, list of the referenced columns</td>
</tr>
<tr>
<td>conpfeqop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If a foreign key, list of the equality operators for (PK = FK) comparisons</td>
</tr>
<tr>
<td>conppeqop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If a foreign key, list of the equality operators for (PK = PK) comparisons</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>conffeqop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If a foreign key, list of the equality operators for FK = FK comparisons</td>
</tr>
<tr>
<td>conexclop</td>
<td>oid[]</td>
<td>pg_operator.oid</td>
<td>If an exclusion constraint, list of the per-column exclusion operators</td>
</tr>
<tr>
<td>conbin</td>
<td>pg_node_tree</td>
<td></td>
<td>If a check constraint, an internal representation of the expression</td>
</tr>
<tr>
<td>consrc</td>
<td>text</td>
<td></td>
<td>If a check constraint, a human-readable representation of the expression</td>
</tr>
</tbody>
</table>

In the case of an exclusion constraint, conkey is only useful for constraint elements that are simple column references. For other cases, a zero appears in conkey and the associated index must be consulted to discover the expression that is constrained. (conkey thus has the same contents as pg_index.indkey for the index.)

**Note:** consrc is not updated when referenced objects change; for example, it won’t track renaming of columns. Rather than relying on this field, it’s best to use pg_get_constraintdef() to extract the definition of a check constraint.

**Note:** pg_class.relchecks needs to agree with the number of check-constraint entries found in this table for each relation.

50.14. pg_conversion

The catalog pg_conversion describes encoding conversion procedures. See CREATE CONVERSION for more information.

Table 50-14. pg_conversion Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>conname</td>
<td>name</td>
<td></td>
<td>Conversion name (unique within a namespace)</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>connamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this conversion</td>
</tr>
<tr>
<td>conowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the conversion</td>
</tr>
<tr>
<td>conforencoding</td>
<td>int4</td>
<td></td>
<td>Source encoding ID</td>
</tr>
<tr>
<td>contoencoding</td>
<td>int4</td>
<td></td>
<td>Destination encoding ID</td>
</tr>
<tr>
<td>conproc</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Conversion procedure</td>
</tr>
<tr>
<td>condefault</td>
<td>bool</td>
<td></td>
<td>True if this is the default conversion</td>
</tr>
</tbody>
</table>

50.15. pg_database

The catalog pg_database stores information about the available databases. Databases are created with the CREATE DATABASE command. Consult Chapter 22 for details about the meaning of some of the parameters.

Unlike most system catalogs, pg_database is shared across all databases of a cluster: there is only one copy of pg_database per cluster, not one per database.

Table 50-15. pg_database Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>datname</td>
<td>name</td>
<td></td>
<td>Database name</td>
</tr>
<tr>
<td>datdba</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the database, usually the user who created it</td>
</tr>
<tr>
<td>encoding</td>
<td>int4</td>
<td></td>
<td>Character encoding for this database (pg_encoding_to_char() can translate this number to the encoding name)</td>
</tr>
<tr>
<td>datcollate</td>
<td>name</td>
<td></td>
<td>LC_COLLATE for this database</td>
</tr>
<tr>
<td>datctype</td>
<td>name</td>
<td></td>
<td>LC_CTYPE for this database</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>--------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>datistemplate</td>
<td>bool</td>
<td></td>
<td>If true, then this database can be cloned by any user with \texttt{CREATEDB} privileges; if false, then only superusers or the owner of the database can clone it.</td>
</tr>
<tr>
<td>datallowconn</td>
<td>bool</td>
<td></td>
<td>If false then no one can connect to this database. This is used to protect the \texttt{template0} database from being altered.</td>
</tr>
<tr>
<td>datconnlimit</td>
<td>int4</td>
<td></td>
<td>Sets maximum number of concurrent connections that can be made to this database. \texttt{-1} means no limit.</td>
</tr>
<tr>
<td>datlastsysoid</td>
<td>oid</td>
<td></td>
<td>Last system OID in the database; useful particularly to \texttt{pg_dump}</td>
</tr>
<tr>
<td>datfrozenxid</td>
<td>xid</td>
<td></td>
<td>All transaction IDs before this one have been replaced with a permanent (&quot;frozen&quot;) transaction ID in this database. This is used to track whether the database needs to be vacuumed in order to prevent transaction ID wraparound or to allow \texttt{pg_clog} to be shrunk. It is the minimum of the per-table \texttt{pg_class.relfrozenxid} values.</td>
</tr>
</tbody>
</table>
### 50.16. **pg_db_role_setting**

The catalog `pg_db_role_setting` records the default values that have been set for run-time configuration variables, for each role and database combination.

Unlike most system catalogs, `pg_db_role_setting` is shared across all databases of a cluster: there is only one copy of `pg_db_role_setting` per cluster, not one per database.

#### Table 50-16. `pg_db_role_setting` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>setdatabase</td>
<td>oid</td>
<td><code>pg_database.oid</code></td>
<td>The OID of the database the setting is applicable to, or zero if not database-specific</td>
</tr>
</tbody>
</table>
### 50.17. `pg_default_acl`

The catalog `pg_default_acl` stores initial privileges to be assigned to newly created objects.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>oid</code></td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td><code>defaclrole</code></td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>The OID of the role associated with this entry</td>
</tr>
<tr>
<td><code>defaclnamespace</code></td>
<td>oid</td>
<td><code>pg_namespace.oid</code></td>
<td>The OID of the namespace associated with this entry, or 0 if none</td>
</tr>
<tr>
<td><code>defaclobjtype</code></td>
<td>char</td>
<td></td>
<td>Type of object this entry is for: (r=) relation (table, view), (S=) sequence, (F=) function, (T=) type</td>
</tr>
<tr>
<td><code>defaclacl</code></td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges that this type of object should have on creation</td>
</tr>
</tbody>
</table>

A `pg_default_acl` entry shows the initial privileges to be assigned to an object belonging to the indicated user. There are currently two types of entry: “global” entries with `defaclnamespace = 0`, and “per-schema” entries that reference a particular schema. If a global entry is present then it overrides the normal hard-wired default privileges for the object type. A per-schema entry, if present, represents privileges to be added to the global or hard-wired default privileges.

Note that when an ACL entry in another catalog is null, it is taken to represent the hard-wired default privileges for its object, not whatever might be in `pg_default_acl` at the moment. `pg_default_acl` is only consulted during object creation.
50.18. pg_depend

The catalog pg_depend records the dependency relationships between database objects. This information allows DROP commands to find which other objects must be dropped by DROP CASCADE or prevent dropping in the DROP RESTRICT case.

See also pg_shdepend, which performs a similar function for dependencies involving objects that are shared across a database cluster.

Table 50-18. pg_depend Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>classid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the dependent object is in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific dependent object</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a table column, this is the column number (the objid and classid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>refclassid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the referenced object is in</td>
</tr>
<tr>
<td>refobjid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific referenced object</td>
</tr>
<tr>
<td>refobjsubid</td>
<td>int4</td>
<td></td>
<td>For a table column, this is the column number (the refobjid and refclassid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>deptype</td>
<td>char</td>
<td></td>
<td>A code defining the specific semantics of this dependency relationship; see text</td>
</tr>
</tbody>
</table>

In all cases, a pg_depend entry indicates that the referenced object cannot be dropped without also dropping the dependent object. However, there are several subflavors identified by deptype:

DEPENDENCY_NORMAL (n)

A normal relationship between separately-created objects. The dependent object can be dropped without affecting the referenced object. The referenced object can only be dropped by specifying CASCADE, in which case the dependent object is dropped, too. Example: a table column has a normal dependency on its data type.
Chapter 50. System Catalogs

**DEPENDENCY_AUTO (a)**

The dependent object can be dropped separately from the referenced object, and should be automatically dropped (regardless of RESTRICT or CASCADE mode) if the referenced object is dropped. Example: a named constraint on a table is made autodependent on the table, so that it will go away if the table is dropped.

**DEPENDENCY_INTERNAL (i)**

The dependent object was created as part of creation of the referenced object, and is really just a part of its internal implementation. A DROP of the dependent object will be disallowed outright (we’ll tell the user to issue a DROP against the referenced object, instead). A DROP of the referenced object will be propagated through to drop the dependent object whether CASCADE is specified or not. Example: a trigger that’s created to enforce a foreign-key constraint is made internally dependent on the constraint’s `pg_constraint` entry.

**DEPENDENCY_EXTENSION (e)**

The dependent object is a member of the extension that is the referenced object (see `pg_extension`). The dependent object can be dropped only via DROP EXTENSION on the referenced object. Functionally this dependency type acts the same as an internal dependency, but it’s kept separate for clarity and to simplify pg_dump.

**DEPENDENCY_AUTO_EXTENSION (x)**

The dependent object is not a member of the extension that is the referenced object (and so should not be ignored by pg_dump), but cannot function without it and should be dropped when the extension itself is. The dependent object may be dropped on its own as well.

**DEPENDENCY_PIN (p)**

There is no dependent object; this type of entry is a signal that the system itself depends on the referenced object, and so that object must never be deleted. Entries of this type are created only by initdb. The columns for the dependent object contain zeroes.

Other dependency flavors might be needed in future.

### 50.19. `pg_description`

The catalog `pg_description` stores optional descriptions (comments) for each database object. Descriptions can be manipulated with the COMMENT command and viewed with psql’s \d commands. Descriptions of many built-in system objects are provided in the initial contents of `pg_description`.

See also `pg_shdescription`, which performs a similar function for descriptions involving objects that are shared across a database cluster.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>objoid</code></td>
<td><code>oid</code></td>
<td>any OID column</td>
<td>The OID of the object this description pertains to</td>
</tr>
<tr>
<td><code>classoid</code></td>
<td><code>oid</code></td>
<td><code>pg_class.oid</code></td>
<td>The OID of the system catalog this object appears in</td>
</tr>
</tbody>
</table>
50.20. \texttt{pg\_enum}

The \texttt{pg\_enum} catalog contains entries showing the values and labels for each enum type. The internal representation of a given enum value is actually the OID of its associated row in \texttt{pg\_enum}.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>enumtypid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>The OID of the \texttt{pg_type} entry owning this enum value</td>
</tr>
<tr>
<td>enumsortorder</td>
<td>float4</td>
<td></td>
<td>The sort position of this enum value within its enum type</td>
</tr>
<tr>
<td>enumlabel</td>
<td>name</td>
<td></td>
<td>The textual label for this enum value</td>
</tr>
</tbody>
</table>

The OIDs for \texttt{pg\_enum} rows follow a special rule: even-numbered OIDs are guaranteed to be ordered in the same way as the sort ordering of their enum type. That is, if two even OIDs belong to the same enum type, the smaller OID must have the smaller \texttt{enumsortorder} value. Odd-numbered OID values need bear no relationship to the sort order. This rule allows the enum comparison routines to avoid catalog lookups in many common cases. The routines that create and alter enum types attempt to assign even OIDs to enum values whenever possible.

When an enum type is created, its members are assigned sort-order positions 1..\texttt{n}. But members added later might be given negative or fractional values of \texttt{enumsortorder}. The only requirement on these values is that they be correctly ordered and unique within each enum type.

50.21. \texttt{pg\_event\_trigger}

The catalog \texttt{pg\_event\_trigger} stores event triggers. See Chapter 38 for more information.
Table 50-21. pg_event_trigger Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>evtname</td>
<td>name</td>
<td></td>
<td>Trigger name (must be unique)</td>
</tr>
<tr>
<td>evtevent</td>
<td>name</td>
<td></td>
<td>Identifies the event for which this trigger fires</td>
</tr>
<tr>
<td>evtowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the event trigger</td>
</tr>
<tr>
<td>evtfoid</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>The function to be called</td>
</tr>
<tr>
<td>evtenabled</td>
<td>char</td>
<td></td>
<td>Controls in which session_replication_role modes the event trigger fires. O = trigger fires in “origin” and “local” modes, D = trigger is disabled, R = trigger fires in “replica” mode, A = trigger fires always.</td>
</tr>
<tr>
<td>evttags</td>
<td>text[]</td>
<td></td>
<td>Command tags for which this trigger will fire. If NULL, the firing of this trigger is not restricted on the basis of the command tag.</td>
</tr>
</tbody>
</table>

50.22. pg_extension

The catalog pg_extension stores information about the installed extensions. See Section 36.15 for details about extensions.

Table 50-22. pg_extension Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>extname</td>
<td>name</td>
<td></td>
<td>Name of the extension</td>
</tr>
<tr>
<td>extowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the extension</td>
</tr>
<tr>
<td>extnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>Schema containing the extension’s exported objects</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-------------------</td>
<td>----------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>extrelocatable</td>
<td>bool</td>
<td></td>
<td>True if extension can be relocated to another schema</td>
</tr>
<tr>
<td>extversion</td>
<td>text</td>
<td></td>
<td>Version name for the extension</td>
</tr>
<tr>
<td>extconfig</td>
<td>oid[]</td>
<td>pg_class.oid</td>
<td>Array of regclass OIDs for the extension’s configuration table(s), or NULL if none</td>
</tr>
<tr>
<td>extcondition</td>
<td>text[]</td>
<td></td>
<td>Array of WHERE-clause filter conditions for the extension’s configuration table(s), or NULL if none</td>
</tr>
</tbody>
</table>

Note that unlike most catalogs with a “namespace” column, extnamespace is not meant to imply that the extension belongs to that schema. Extension names are never schema-qualified. Rather, extnamespace indicates the schema that contains most or all of the extension’s objects. If extrelocatable is true, then this schema must in fact contain all schema-qualifiable objects belonging to the extension.

### 50.23. pg_foreign_data_wrapper

The catalog pg_foreign_data_wrapper stores foreign-data wrapper definitions. A foreign-data wrapper is the mechanism by which external data, residing on foreign servers, is accessed.

Table 50-23. pg_foreign_data_wrapper Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>fdwname</td>
<td>name</td>
<td></td>
<td>Name of the foreign-data wrapper</td>
</tr>
<tr>
<td>fdwowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the foreign-data wrapper</td>
</tr>
<tr>
<td>fdwhandler</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>References a handler function that is responsible for supplying execution routines for the foreign-data wrapper. Zero if no handler is provided</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>fdwvalidator</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>References a validator function that is responsible for checking the validity of the options given to the foreign-data wrapper, as well as options for foreign servers and user mappings using the foreign-data wrapper. Zero if no validator is provided</td>
</tr>
<tr>
<td>fdwacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
<tr>
<td>fdwoptions</td>
<td>text[]</td>
<td></td>
<td>Foreign-data wrapper specific options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

#### 50.24. pg_foreign_server

The catalog `pg_foreign_server` stores foreign server definitions. A foreign server describes a source of external data, such as a remote server. Foreign servers are accessed via foreign-data wrappers.

**Table 50-24. pg_foreign_server Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>srvname</td>
<td>name</td>
<td></td>
<td>Name of the foreign server</td>
</tr>
<tr>
<td>srvowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the foreign server</td>
</tr>
<tr>
<td>srvfdw</td>
<td>oid</td>
<td>pg_foreign_data_wrapper.oid</td>
<td>OID of the foreign-data wrapper of this foreign server</td>
</tr>
<tr>
<td>srvtype</td>
<td>text</td>
<td></td>
<td>Type of the server (optional)</td>
</tr>
<tr>
<td>srvversion</td>
<td>text</td>
<td></td>
<td>Version of the server (optional)</td>
</tr>
<tr>
<td>srvacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>srvoptions</td>
<td>text[]</td>
<td></td>
<td>Foreign server specific options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

### 50.25. pg_foreign_table

The catalog `pg_foreign_table` contains auxiliary information about foreign tables. A foreign table is primarily represented by a `pg_class` entry, just like a regular table. Its `pg_foreign_table` entry contains the information that is pertinent only to foreign tables and not any other kind of relation.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ftrelid</td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>OID of the <code>pg_class</code> entry for this foreign table</td>
</tr>
<tr>
<td>ftserver</td>
<td>oid</td>
<td><code>pg_foreign_server.oid</code></td>
<td>OID of the foreign server for this foreign table</td>
</tr>
<tr>
<td>ftoptions</td>
<td>text[]</td>
<td></td>
<td>Foreign table options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

### 50.26. pg_index

The catalog `pg_index` contains part of the information about indexes. The rest is mostly in `pg_class`.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>indexrelid</td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>The OID of the <code>pg_class</code> entry for this index</td>
</tr>
<tr>
<td>indrelid</td>
<td>oid</td>
<td><code>pg_class.oid</code></td>
<td>The OID of the <code>pg_class</code> entry for the table this index is for</td>
</tr>
<tr>
<td>indnatts</td>
<td>int2</td>
<td></td>
<td>The number of columns in the index (duplicates <code>pg_class.relnatts</code>)</td>
</tr>
<tr>
<td>indisunique</td>
<td>bool</td>
<td></td>
<td>If true, this is a unique index</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>-------</td>
<td>------------</td>
<td>------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>indisprimary</td>
<td>bool</td>
<td></td>
<td>If true, this index represents the primary key of the table (indisunique should always be true when this is true)</td>
</tr>
<tr>
<td>indisexclusion</td>
<td>bool</td>
<td></td>
<td>If true, this index supports an exclusion constraint</td>
</tr>
<tr>
<td>indimmediate</td>
<td>bool</td>
<td></td>
<td>If true, the uniqueness check is enforced immediately on insertion (irrelevant if indisunique is not true)</td>
</tr>
<tr>
<td>indisclustered</td>
<td>bool</td>
<td></td>
<td>If true, the table was last clustered on this index</td>
</tr>
<tr>
<td>indisvalid</td>
<td>bool</td>
<td></td>
<td>If true, the index is currently valid for queries. False means the index is possibly incomplete: it must still be modified by INSERT/UPDATE operations, but it cannot safely be used for queries. If it is unique, the uniqueness property is not guaranteed true either.</td>
</tr>
<tr>
<td>indcheckxmin</td>
<td>bool</td>
<td></td>
<td>If true, queries must not use the index until the xmin of this pg_index row is below their TransactionXmin event horizon, because the table may contain broken HOT chains with incompatible rows that they can see</td>
</tr>
<tr>
<td>indisready</td>
<td>bool</td>
<td></td>
<td>If true, the index is currently ready for inserts. False means the index must be ignored by INSERT/UPDATE operations.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>----------------</td>
<td>----------------</td>
<td>-----------------------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>indislive</td>
<td>bool</td>
<td></td>
<td>If false, the index is in process of being dropped, and should be ignored for all purposes (including HOT-safety decisions)</td>
</tr>
<tr>
<td>indisreplident</td>
<td>bool</td>
<td></td>
<td>If true this index has been chosen as “replica identity” using ALTER TABLE ... REPLICA IDENTITY USING INDEX ...</td>
</tr>
<tr>
<td>indkey</td>
<td>int2vector</td>
<td>pg_attribute.attnum</td>
<td>This is an array of indnatts values that indicate which table columns this index indexes. For example a value of 1 3 would mean that the first and the third table columns make up the index key. A zero in this array indicates that the corresponding index attribute is an expression over the table columns, rather than a simple column reference.</td>
</tr>
<tr>
<td>indcollation</td>
<td>oidvector</td>
<td>pg_collation.oid</td>
<td>For each column in the index key, this contains the OID of the collation to use for the index.</td>
</tr>
<tr>
<td>indclass</td>
<td>oidvector</td>
<td>pg_opclass.oid</td>
<td>For each column in the index key, this contains the OID of the operator class to use. See pg_opclass for details.</td>
</tr>
<tr>
<td>indoption</td>
<td>int2vector</td>
<td></td>
<td>This is an array of indnatts values that store per-column flag bits. The meaning of the bits is defined by the index’s access method.</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>indexprs</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression trees (in nodeToString() representation) for index attributes that are not simple column references. This is a list with one element for each zero entry in indkey. Null if all index attributes are simple references.</td>
</tr>
<tr>
<td>indpred</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression tree (in nodeToString() representation) for partial index predicate. Null if not a partial index.</td>
</tr>
</tbody>
</table>

50.27. pg_inherits

The catalog pg_inherits records information about table inheritance hierarchies. There is one entry for each direct child table in the database. (Indirect inheritance can be determined by following chains of entries.)

Table 50-27. pg_inherits Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>inhrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the child table</td>
</tr>
<tr>
<td>inhparent</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the parent table</td>
</tr>
<tr>
<td>inhseqno</td>
<td>int4</td>
<td></td>
<td>If there is more than one direct parent for a child table (multiple inheritance), this number tells the order in which the inherited columns are to be arranged. The count starts at 1.</td>
</tr>
</tbody>
</table>

50.28. pg_init_privs

The catalog pg_init_privs records information about the initial privileges of objects in the system. There is one entry for each object in the database which has a non-default (non-NULL) initial set of
privileges.

Objects can have initial privileges either by having those privileges set when the system is initialized (by initdb) or when the object is created during a `CREATE EXTENSION` and the extension script sets initial privileges using the `GRANT` system. Note that the system will automatically handle recording of the privileges during the extension script and that extension authors need only use the `GRANT` and `REVOKE` statements in their script to have the privileges recorded. The `privtype` column indicates if the initial privilege was set by initdb or during a `CREATE EXTENSION` command.

Objects which have initial privileges set by initdb will have entries where `privtype` is ‘i’, while objects which have initial privileges set by `CREATE EXTENSION` will have entries where `privtype` is ‘e’.

### Table 50-28. `pg_init_privs` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific object</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td><code>pg_class</code>.oid</td>
<td>The OID of the system catalog the object is in</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a table column, this is the column number (the <code>objoid</code> and <code>classoid</code> refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>privtype</td>
<td>char</td>
<td></td>
<td>A code defining the type of initial privilege of this object; see text</td>
</tr>
<tr>
<td>initprivs</td>
<td>aclitem[]</td>
<td></td>
<td>The initial access privileges; see <code>GRANT</code> and <code>REVOKE</code> for details</td>
</tr>
</tbody>
</table>

### 50.29. `pg_language`

The catalog `pg_language` registers languages in which you can write functions or stored procedures. See `CREATE LANGUAGE` and Chapter 40 for more information about language handlers.

### Table 50-29. `pg_language` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>lanname</td>
<td>name</td>
<td></td>
<td>Name of the language</td>
</tr>
<tr>
<td>lanowner</td>
<td>oid</td>
<td><code>pg_authid</code>.oid</td>
<td>Owner of the language</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>-------</td>
<td>------------</td>
<td>--------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>lanispl</td>
<td>bool</td>
<td></td>
<td>This is false for internal languages (such as SQL) and true for user-defined languages. Currently, pg_dump still uses this to determine which languages need to be dumped, but this might be replaced by a different mechanism in the future.</td>
</tr>
<tr>
<td>lanpltrusted</td>
<td>bool</td>
<td></td>
<td>True if this is a trusted language, which means that it is believed not to grant access to anything outside the normal SQL execution environment. Only superusers can create functions in untrusted languages.</td>
</tr>
<tr>
<td>lanplcallfoid</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>For noninternal languages this references the language handler, which is a special function that is responsible for executing all functions that are written in the particular language.</td>
</tr>
<tr>
<td>laninline</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>This references a function that is responsible for executing “inline” anonymous code blocks (DO blocks). Zero if inline blocks are not supported.</td>
</tr>
<tr>
<td>lanvalidator</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>This references a language validator function that is responsible for checking the syntax and validity of new functions when they are created. Zero if no validator is provided.</td>
</tr>
</tbody>
</table>
### 50.30. \texttt{pg\_largeobject}

The catalog \texttt{pg\_largeobject} holds the data making up “large objects”. A large object is identified by an OID assigned when it is created. Each large object is broken into segments or “pages” small enough to be conveniently stored as rows in \texttt{pg\_largeobject}. The amount of data per page is defined to be \texttt{LOBLKSIZE} (which is currently \texttt{BLCKSZ/4}, or typically 2 kB).

Prior to PostgreSQL 9.0, there was no permission structure associated with large objects. As a result, \texttt{pg\_largeobject} was publicly readable and could be used to obtain the OIDs (and contents) of all large objects in the system. This is no longer the case; use \texttt{pg\_largeobject\_metadata} to obtain a list of large object OIDs.

#### Table 50-30. \texttt{pg\_largeobject} Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{loid}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_largeobject_metadata.oid}</td>
<td>Identifier of the large object that includes this page</td>
</tr>
<tr>
<td>\texttt{pageno}</td>
<td>\texttt{int4}</td>
<td></td>
<td>Page number of this page within its large object (counting from zero)</td>
</tr>
<tr>
<td>\texttt{data}</td>
<td>\texttt{bytea}</td>
<td></td>
<td>Actual data stored in the large object. This will never be more than \texttt{LOBLKSIZE} bytes and might be less.</td>
</tr>
</tbody>
</table>

Each row of \texttt{pg\_largeobject} holds data for one page of a large object, beginning at byte offset \((\text{pageno} \times \text{LOBLKSIZE})\) within the object. The implementation allows sparse storage: pages might be missing, and might be shorter than \texttt{LOBLKSIZE} bytes even if they are not the last page of the object. Missing regions within a large object read as zeroes.

### 50.31. \texttt{pg\_largeobject\_metadata}

The catalog \texttt{pg\_largeobject\_metadata} holds metadata associated with large objects. The actual large object data is stored in \texttt{pg\_largeobject}.

#### Table 50-31. \texttt{pg\_largeobject\_metadata} Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

1987
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>lomowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the large object</td>
</tr>
<tr>
<td>lomacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
</tbody>
</table>

### 50.32. pg_namespace

The catalog pg_namespace stores namespaces. A namespace is the structure underlying SQL schemas: each namespace can have a separate collection of relations, types, etc. without name conflicts.

Table 50-32. pg_namespace Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>nspname</td>
<td>name</td>
<td></td>
<td>Name of the namespace</td>
</tr>
<tr>
<td>nspowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the namespace</td>
</tr>
<tr>
<td>nspacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
</tbody>
</table>

### 50.33. pg_opclass

The catalog pg_opclass defines index access method operator classes. Each operator class defines semantics for index columns of a particular data type and a particular index access method. An operator class essentially specifies that a particular operator family is applicable to a particular indexable column data type. The set of operators from the family that are actually usable with the indexed column are whichever ones accept the column’s data type as their left-hand input.

Operator classes are described at length in Section 36.14.

Table 50-33. pg_opclass Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>opcmethod</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>Index access method operator class is for</td>
</tr>
<tr>
<td>opcname</td>
<td>name</td>
<td></td>
<td>Name of this operator class</td>
</tr>
<tr>
<td>opcnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>Namespace of this operator class</td>
</tr>
<tr>
<td>opcowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the operator class</td>
</tr>
<tr>
<td>opcfamily</td>
<td>oid</td>
<td>pg_opfamily.oid</td>
<td>Operator family containing the operator class</td>
</tr>
<tr>
<td>opcintype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type that the operator class indexes</td>
</tr>
<tr>
<td>opcdefault</td>
<td>bool</td>
<td></td>
<td>True if this operator class is the default for</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>opcintype</td>
</tr>
<tr>
<td>opckeytype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of data stored in index, or zero if same as</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>opcintype</td>
</tr>
</tbody>
</table>

An operator class’s opcmethod must match the opfmethod of its containing operator family. Also, there must be no more than one pg_opclass row having opcdefault true for any given combination of opcmethod and opcintype.

### 50.34. pg_operator

The catalog pg_operator stores information about operators. See CREATE OPERATOR and Section 36.12 for more information.

Table 50-34. pg_operator Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>oprname</td>
<td>name</td>
<td></td>
<td>Name of the operator</td>
</tr>
<tr>
<td>oprnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this operator</td>
</tr>
<tr>
<td>oprowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the operator</td>
</tr>
<tr>
<td>oprkind</td>
<td>char</td>
<td></td>
<td>b = infix (“both”), l = prefix (“left”), r = postfix (“right”)</td>
</tr>
<tr>
<td>oprcanmerge</td>
<td>bool</td>
<td></td>
<td>This operator supports merge joins</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oprcanhash</td>
<td>bool</td>
<td>pg_type</td>
<td>This operator supports hash joins</td>
</tr>
<tr>
<td>oprleft</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of the left operand</td>
</tr>
<tr>
<td>oprright</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of the right operand</td>
</tr>
<tr>
<td>oprresult</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Type of the result</td>
</tr>
<tr>
<td>oprcom</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>Commutator of this operator, if any</td>
</tr>
<tr>
<td>oprnegate</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>Negator of this operator, if any</td>
</tr>
<tr>
<td>oprcode</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Function that implements this operator</td>
</tr>
<tr>
<td>oprrest</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Restriction selectivity estimation function for this operator</td>
</tr>
<tr>
<td>oprjoin</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Join selectivity estimation function for this operator</td>
</tr>
</tbody>
</table>

Unused column contain zeroes. For example, oprleft is zero for a prefix operator.

50.35. pg_opfamily

The catalog pg_opfamily defines operator families. Each operator family is a collection of operators and associated support routines that implement the semantics specified for a particular index access method. Furthermore, the operators in a family are all "compatible", in a way that is specified by the access method. The operator family concept allows cross-data-type operators to be used with indexes and to be reasoned about using knowledge of access method semantics.

Operator families are described at length in Section 36.14.

Table 50-35. pg_opfamily Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>opfmethod</td>
<td>oid</td>
<td>pg_am.oid</td>
<td>Index access method operator family is for</td>
</tr>
<tr>
<td>opfname</td>
<td>name</td>
<td></td>
<td>Name of this operator family</td>
</tr>
<tr>
<td>opfnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>Namespace of this operator family</td>
</tr>
<tr>
<td>opfowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the operator family</td>
</tr>
</tbody>
</table>
The majority of the information defining an operator family is not in its `pg_opfamily` row, but in the associated rows in `pg_amop`, `pg_amproc`, and `pg_opclass`.

### 50.36. `pg_pltemplate`

The catalog `pg_pltemplate` stores “template” information for procedural languages. A template for a language allows the language to be created in a particular database by a simple `CREATE LANGUAGE` command, with no need to specify implementation details.

Unlike most system catalogs, `pg_pltemplate` is shared across all databases of a cluster: there is only one copy of `pg_pltemplate` per cluster, not one per database. This allows the information to be accessible in each database as it is needed.

#### Table 50-36. `pg_pltemplate` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>tmplname</td>
<td>name</td>
<td>Name of the language this template is for</td>
</tr>
<tr>
<td>tmpltrusted</td>
<td>boolean</td>
<td>True if language is considered trusted</td>
</tr>
<tr>
<td>tmpldbcreate</td>
<td>boolean</td>
<td>True if language may be created by a database owner</td>
</tr>
<tr>
<td>tmplhandler</td>
<td>text</td>
<td>Name of call handler function</td>
</tr>
<tr>
<td>tmplinline</td>
<td>text</td>
<td>Name of anonymous-block handler function, or null if none</td>
</tr>
<tr>
<td>tmplvalidator</td>
<td>text</td>
<td>Name of validator function, or null if none</td>
</tr>
<tr>
<td>tmpllibrary</td>
<td>text</td>
<td>Path of shared library that implements language</td>
</tr>
<tr>
<td>tmplacl</td>
<td>aclitem[]</td>
<td>Access privileges for template (not actually used)</td>
</tr>
</tbody>
</table>

There are not currently any commands that manipulate procedural language templates; to change the built-in information, a superuser must modify the table using ordinary `INSERT`, `DELETE`, or `UPDATE` commands.

**Note:** It is likely that `pg_pltemplate` will be removed in some future release of PostgreSQL, in favor of keeping this knowledge about procedural languages in their respective extension installation scripts.

### 50.37. `pg_policy`

The catalog `pg_policy` stores row level security policies for tables. A policy includes the kind of command that it applies to (possibly all commands), the roles that it applies to, the expression to be
added as a security-barrier qualification to queries that include the table, and the expression to be added as a WITH CHECK option for queries that attempt to add new records to the table.

Table 50-37. pg_policy Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>polname</td>
<td>name</td>
<td></td>
<td>The name of the policy</td>
</tr>
<tr>
<td>polrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table to which the policy applies</td>
</tr>
<tr>
<td>polcmd</td>
<td>char</td>
<td></td>
<td>The command type to which the policy is applied: r for SELECT, a for INSERT, w for UPDATE, d for DELETE, or * for all</td>
</tr>
<tr>
<td>polroles</td>
<td>oid[]</td>
<td>pg_authid.oid</td>
<td>The roles to which the policy is applied</td>
</tr>
<tr>
<td>polqual</td>
<td>pg_node_tree</td>
<td></td>
<td>The expression tree to be added to the security barrier qualifications for queries that use the table</td>
</tr>
<tr>
<td>polwithcheck</td>
<td>pg_node_tree</td>
<td></td>
<td>The expression tree to be added to the WITH CHECK qualifications for queries that attempt to add rows to the table</td>
</tr>
</tbody>
</table>

**Note:** Policies stored in pg_policy are applied only when pg_class.relrowsecurity is set for their table.

50.38. pg_proc

The catalog pg_proc stores information about functions (or procedures). See CREATE FUNCTION and Section 36.3 for more information.

The table contains data for aggregate functions as well as plain functions. If proisagg is true, there should be a matching row in pg_aggregate.

Table 50-38. pg_proc Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>proname</td>
<td>name</td>
<td></td>
<td>Name of the function</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>--------</td>
<td>------------</td>
<td>------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>pronamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this function</td>
</tr>
<tr>
<td>proowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the function</td>
</tr>
<tr>
<td>prolang</td>
<td>oid</td>
<td>pg_language.oid</td>
<td>Implementation language or call interface of this function</td>
</tr>
<tr>
<td>procost</td>
<td>float4</td>
<td></td>
<td>Estimated execution cost (in units of cpu_operator_cost); if proretset, this is cost per row returned</td>
</tr>
<tr>
<td>prorows</td>
<td>float4</td>
<td></td>
<td>Estimated number of result rows (zero if not proretset)</td>
</tr>
<tr>
<td>provariadic</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>Data type of the variadic array parameter’s elements, or zero if the function does not have a variadic parameter</td>
</tr>
<tr>
<td>protransform</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Calls to this function can be simplified by this other function (see Section 36.9.11)</td>
</tr>
<tr>
<td>proisagg</td>
<td>bool</td>
<td></td>
<td>Function is an aggregate function</td>
</tr>
<tr>
<td>proiswindow</td>
<td>bool</td>
<td></td>
<td>Function is a window function</td>
</tr>
<tr>
<td>prosecdef</td>
<td>bool</td>
<td></td>
<td>Function is a security definer (i.e., a “setuid” function)</td>
</tr>
<tr>
<td>proleakproof</td>
<td>bool</td>
<td></td>
<td>The function has no side effects. No information about the arguments is conveyed except via the return value. Any function that might throw an error depending on the values of its arguments is not leak-proof.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-------------</td>
<td>--------</td>
<td>------------</td>
<td>--------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>proisstrict</td>
<td>bool</td>
<td></td>
<td>Function returns null if any call argument is null. In that case the function won’t actually be called at all. Functions that are not “strict” must be prepared to handle null inputs.</td>
</tr>
<tr>
<td>proretset</td>
<td>bool</td>
<td></td>
<td>Function returns a set (i.e., multiple values of the specified data type)</td>
</tr>
<tr>
<td>provolatile</td>
<td>char</td>
<td></td>
<td>provolatile tells whether the function’s result depends only on its input arguments, or is affected by outside factors. It is i for “immutable” functions, which always deliver the same result for the same inputs. It is s for “stable” functions, whose results (for fixed inputs) do not change within a scan. It is v for “volatile” functions, whose results might change at any time. (Use v also for functions with side-effects, so that calls to them cannot get optimized away.)</td>
</tr>
</tbody>
</table>
### proparallel

- **Name**: proparallel
- **Type**: char
- **References**: proparallel
- **Description**: proparallel tells whether the function can be safely run in parallel mode. It is `s` for functions which are safe to run in parallel mode without restriction. It is `r` for functions which can be run in parallel mode, but their execution is restricted to the parallel group leader; parallel worker processes cannot invoke these functions. It is `u` for functions which are unsafe in parallel mode; the presence of such a function forces a serial execution plan.

### pronargs

- **Name**: pronargs
- **Type**: int2
- **References**: Number of input arguments
- **Description**:

### pronargdefaults

- **Name**: pronargdefaults
- **Type**: int2
- **References**: Number of arguments that have defaults
- **Description**: A function that indicates whether a function has one or more arguments that have default values.

### prorettype

- **Name**: prorettype
- **Type**: oid
- **References**: pg_type.oid
- **Description**: Data type of the return value

### proargtypes

- **Name**: proargtypes
- **Type**: oidvector
- **References**: pg_type.oid
- **Description**: An array with the data types of the function arguments. This includes only input arguments (including INOUT and VARIADIC arguments), and thus represents the call signature of the function.
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>proallargtypes</td>
<td>oid[]</td>
<td>pg_type.oid</td>
<td>An array with the data types of the function arguments. This includes all arguments (including <code>OUT</code> and <code>INOUT</code> arguments); however, if all the arguments are <code>IN</code> arguments, this field will be null. Note that subscripting is 1-based, whereas for historical reasons <code>proargtypes</code> is subscripted from 0.</td>
</tr>
<tr>
<td>proargmodes</td>
<td>char[]</td>
<td></td>
<td>An array with the modes of the function arguments, encoded as <code>i</code> for <code>IN</code> arguments, <code>o</code> for <code>OUT</code> arguments, <code>b</code> for <code>INOUT</code> arguments, <code>v</code> for <code>VARIADIC</code> arguments, <code>t</code> for <code>TABLE</code> arguments. If all the arguments are <code>IN</code> arguments, this field will be null. Note that subscripts correspond to positions of <code>proallargtypes</code> not <code>proargtypes</code>.</td>
</tr>
<tr>
<td>proargnames</td>
<td>text[]</td>
<td></td>
<td>An array with the names of the function arguments. Arguments without a name are set to empty strings in the array. If none of the arguments have a name, this field will be null. Note that subscripts correspond to positions of <code>proallargtypes</code> not <code>proargtypes</code>.</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>proargdefaults</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression trees (in nodeToString() representation) for default values. This is a list with proargdefaults elements, corresponding to the last N input arguments (i.e., the last N proargtypes positions). If none of the arguments have defaults, this field will be null.</td>
</tr>
<tr>
<td>protrftypes</td>
<td>oid[]</td>
<td></td>
<td>Data type OIDs for which to apply transforms.</td>
</tr>
<tr>
<td>prosrc</td>
<td>text</td>
<td></td>
<td>This tells the function handler how to invoke the function. It might be the actual source code of the function for interpreted languages, a link symbol, a file name, or just about anything else, depending on the implementation language/call convention.</td>
</tr>
<tr>
<td>probin</td>
<td>text</td>
<td></td>
<td>Additional information about how to invoke the function. Again, the interpretation is language-specific.</td>
</tr>
<tr>
<td>proconfig</td>
<td>text[]</td>
<td></td>
<td>Function’s local settings for run-time configuration variables.</td>
</tr>
<tr>
<td>proacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details.</td>
</tr>
</tbody>
</table>

For compiled functions, both built-in and dynamically loaded, prosrc contains the function’s C-language name (link symbol). For all other currently-known language types, prosrc contains the function’s source text. probin is unused except for dynamically-loaded C functions, for which it gives the name of the shared library file containing the function.
50.39. pg_range

The catalog pg_range stores information about range types. This is in addition to the types’ entries in pg_type.

Table 50-39. pg_range Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rngtypid</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the range type</td>
</tr>
<tr>
<td>rngsubtype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the element type (subtype) of this range type</td>
</tr>
<tr>
<td>rngcollation</td>
<td>oid</td>
<td>pg_collation.oid</td>
<td>OID of the collation used for range comparisons, or 0 if none</td>
</tr>
<tr>
<td>rngsubopc</td>
<td>oid</td>
<td>pg_opclass.oid</td>
<td>OID of the subtype’s operator class used for range comparisons</td>
</tr>
<tr>
<td>rngcanonical</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the function to convert a range value into canonical form, or 0 if none</td>
</tr>
<tr>
<td>rngsubdiff</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the function to return the difference between two element values as double precision, or 0 if none</td>
</tr>
</tbody>
</table>

rngsubopc (plus rngcollation, if the element type is collatable) determines the sort ordering used by the range type. rngcanonical is used when the element type is discrete. rngsubdiff is optional but should be supplied to improve performance of GiST indexes on the range type.

50.40. pg_replication_origin

The pg_replication_origin catalog contains all replication origins created. For more on replication origins see Chapter 48.

Table 50-40. pg_replication_origin Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>roident</td>
<td>Oid</td>
<td></td>
<td>A unique, cluster-wide identifier for the replication origin. Should never leave the system.</td>
</tr>
</tbody>
</table>
### 50.41. pg_rewrite

The catalog **pg_rewrite** stores rewrite rules for tables and views.

#### Table 50-41. pg_rewrite Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>rulename</td>
<td>name</td>
<td></td>
<td>Rule name</td>
</tr>
<tr>
<td>ev_class</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this rule is for</td>
</tr>
<tr>
<td>ev_type</td>
<td>char</td>
<td></td>
<td>Event type that the rule is for: 1 = SELECT, 2 = UPDATE, 3 = INSERT, 4 = DELETE</td>
</tr>
<tr>
<td>ev_enabled</td>
<td>char</td>
<td></td>
<td>Controls in which session_replication_role modes the rule fires. O = rule fires in “origin” and “local” modes, D = rule is disabled, R = rule fires in “replica” mode, A = rule fires always.</td>
</tr>
<tr>
<td>is_instead</td>
<td>bool</td>
<td></td>
<td>True if the rule is an INSTEAD rule</td>
</tr>
<tr>
<td>ev_qual</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression tree (in the form of a nodeToString() representation) for the rule’s qualifying condition</td>
</tr>
<tr>
<td>ev_action</td>
<td>pg_node_tree</td>
<td></td>
<td>Query tree (in the form of a nodeToString() representation) for the rule’s action</td>
</tr>
</tbody>
</table>

**Note:** pg_class.relhasrules must be true if a table has any rules in this catalog.
50.42. pg_seclabel

The catalog `pg_seclabel` stores security labels on database objects. Security labels can be manipulated with the SECURITY LABEL command. For an easier way to view security labels, see Section 50.74.

See also `pg_shseclabel`, which performs a similar function for security labels of database objects that are shared across a database cluster.

**Table 50-42. pg_seclabel Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this security label pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a security label on a table column, this is the column number (the objoid and classoid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>provider</td>
<td>text</td>
<td></td>
<td>The label provider associated with this label.</td>
</tr>
<tr>
<td>label</td>
<td>text</td>
<td></td>
<td>The security label applied to this object.</td>
</tr>
</tbody>
</table>

50.43. pg_shdepend

The catalog `pg_shdepend` records the dependency relationships between database objects and shared objects, such as roles. This information allows PostgreSQL to ensure that those objects are unreferenced before attempting to delete them.

See also `pg_depend`, which performs a similar function for dependencies involving objects within a single database.

Unlike most system catalogs, `pg_shdepend` is shared across all databases of a cluster: there is only one copy of `pg_shdepend` per cluster, not one per database.

**Table 50-43. pg_shdepend Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>dbid</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>The OID of the database the dependent object is in, or zero for a shared object</td>
</tr>
<tr>
<td>classid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the dependent object is in</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific dependent object</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a table column, this is the column number (the objid and classid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>refclassid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog the referenced object is in (must be a shared catalog)</td>
</tr>
<tr>
<td>refobjid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the specific referenced object</td>
</tr>
<tr>
<td>deptype</td>
<td>char</td>
<td></td>
<td>A code defining the specific semantics of this dependency relationship; see text</td>
</tr>
</tbody>
</table>

In all cases, a `pg_shdepend` entry indicates that the referenced object cannot be dropped without also dropping the dependent object. However, there are several subflavors identified by `deptype`:

**SHARED_DEPENDENCY_OWNER (o)**

The referenced object (which must be a role) is the owner of the dependent object.

**SHARED_DEPENDENCY_ACL (a)**

The referenced object (which must be a role) is mentioned in the ACL (access control list, i.e., privileges list) of the dependent object. (A SHARED_DEPENDENCY_ACL entry is not made for the owner of the object, since the owner will have a SHARED_DEPENDENCY_OWNER entry anyway.)

**SHARED_DEPENDENCY_POLICY (r)**

The referenced object (which must be a role) is mentioned as the target of a dependent policy object.

**SHARED_DEPENDENCY_PIN (p)**

There is no dependent object; this type of entry is a signal that the system itself depends on the referenced object, and so that object must never be deleted. Entries of this type are created only by `initdb`. The columns for the dependent object contain zeroes.

Other dependency flavors might be needed in future. Note in particular that the current definition only supports roles as referenced objects.
50.44. pg_shdescription

The catalog `pg_shdescription` stores optional descriptions (comments) for shared database objects. Descriptions can be manipulated with the COMMENT command and viewed with psql’s \d commands.

See also `pg_description`, which performs a similar function for descriptions involving objects within a single database.

Unlike most system catalogs, `pg_shdescription` is shared across all databases of a cluster: there is only one copy of `pg_shdescription` per cluster, not one per database.

Table 50-44. pg_shdescription Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this description pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>description</td>
<td>text</td>
<td></td>
<td>Arbitrary text that serves as the description of this object</td>
</tr>
</tbody>
</table>

50.45. pg_shseclabel

The catalog `pg_shseclabel` stores security labels on shared database objects. Security labels can be manipulated with the SECURITY LABEL command. For an easier way to view security labels, see Section 50.74.

See also `pg_seclabel`, which performs a similar function for security labels involving objects within a single database.

Unlike most system catalogs, `pg_shseclabel` is shared across all databases of a cluster: there is only one copy of `pg_shseclabel` per cluster, not one per database.

Table 50-45. pg_shseclabel Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this security label pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>provider</td>
<td>text</td>
<td></td>
<td>The label provider associated with this label</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>label</td>
<td>text</td>
<td></td>
<td>The security label applied to this object.</td>
</tr>
</tbody>
</table>

50.46. pg_statistic

The catalog pg_statistic stores statistical data about the contents of the database. Entries are created by ANALYZE and subsequently used by the query planner. Note that all the statistical data is inherently approximate, even assuming that it is up-to-date.

Normally there is one entry, with stainherit = false, for each table column that has been analyzed. If the table has inheritance children, a second entry with stainherit = true is also created. This row represents the column’s statistics over the inheritance tree, i.e., statistics for the data you’d see with SELECT column FROM table*, whereas the stainherit = false row represents the results of SELECT column FROM ONLY table.

gp_statistic also stores statistical data about the values of index expressions. These are described as if they were actual data columns; in particular, starelid references the index. No entry is made for an ordinary non-expression index column, however, since it would be redundant with the entry for the underlying table column. Currently, entries for index expressions always have stainherit = false.

Since different kinds of statistics might be appropriate for different kinds of data, pg_statistic is designed not to assume very much about what sort of statistics it stores. Only extremely general statistics (such as nullness) are given dedicated columns in pg_statistic. Everything else is stored in “slots”, which are groups of associated columns whose content is identified by a code number in one of the slot’s columns. For more information see src/include/catalog/pg_statistic.h.

gp_statistic should not be readable by the public, since even statistical information about a table’s contents might be considered sensitive. (Example: minimum and maximum values of a salary column might be quite interesting.) pg_stats is a publicly readable view on pg_statistic that only exposes information about those tables that are readable by the current user.

Table 50-46. pg_statistic Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>starelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table or index that the described column belongs to</td>
</tr>
<tr>
<td>staattnum</td>
<td>int2</td>
<td>pg_attribute.attnum</td>
<td>The number of the described column</td>
</tr>
<tr>
<td>stainherit</td>
<td>bool</td>
<td></td>
<td>If true, the stats include inheritance child columns, not just the values in the specified relation</td>
</tr>
<tr>
<td>stanullfrac</td>
<td>float4</td>
<td></td>
<td>The fraction of the column’s entries that are null</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-------------</td>
<td>-----------</td>
<td>------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>stawidth</td>
<td>int4</td>
<td></td>
<td>The average stored width, in bytes, of nonnull entries</td>
</tr>
<tr>
<td>stadistinct</td>
<td>float4</td>
<td></td>
<td>The number of distinct nonnull data values in the column. A value greater than zero is the actual number of distinct values. A value less than zero is the negative of a multiplier for the number of rows in the table; for example, a column in which about 80% of the values are nonnull and each nonnull value appears about twice on average could be represented by <code>stadistinct = -0.4</code>. A zero value means the number of distinct values is unknown.</td>
</tr>
<tr>
<td>stakindN</td>
<td>int2</td>
<td></td>
<td>A code number indicating the kind of statistics stored in the Nth “slot” of the <code>pg_statistic</code> row.</td>
</tr>
<tr>
<td>staopN</td>
<td>oid</td>
<td>pg_operator.oid</td>
<td>An operator used to derive the statistics stored in the Nth “slot”. For example, a histogram slot would show the &lt; operator that defines the sort order of the data.</td>
</tr>
<tr>
<td>stanumbersN</td>
<td>float4[]</td>
<td></td>
<td>Numerical statistics of the appropriate kind for the Nth “slot”, or null if the slot kind does not involve numerical values</td>
</tr>
</tbody>
</table>
**Table 50-47. pg_tablespace Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>spcname</td>
<td>name</td>
<td></td>
<td>Tablespace name</td>
</tr>
<tr>
<td>spcowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the tablespace, usually the user who created it</td>
</tr>
<tr>
<td>spcacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details</td>
</tr>
<tr>
<td>spcoptions</td>
<td>text[]</td>
<td></td>
<td>Tablespace-level options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

**50.47. pg_tablespace**

The catalog `pg_tablespace` stores information about the available tablespaces. Tables can be placed in particular tablespaces to aid administration of disk layout.

Unlike most system catalogs, `pg_tablespace` is shared across all databases of a cluster: there is only one copy of `pg_tablespace` per cluster, not one per database.

**50.48. pg_transform**

The catalog `pg_transform` stores information about transforms, which are a mechanism to adapt data types to procedural languages. See CREATE TRANSFORM for more information.
Table 50-48. pg_transform Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>trftype</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>OID of the data type this transform is for</td>
</tr>
<tr>
<td>trflang</td>
<td>oid</td>
<td>pg_language.oid</td>
<td>OID of the language this transform is for</td>
</tr>
<tr>
<td>trffromsql</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>The OID of the function to use when converting the data type for input to the procedural language (e.g., function parameters). Zero is stored if this operation is not supported.</td>
</tr>
<tr>
<td>trftosql</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>The OID of the function to use when converting output from the procedural language (e.g., return values) to the data type. Zero is stored if this operation is not supported.</td>
</tr>
</tbody>
</table>

50.49. pg_trigger

The catalog pg_trigger stores triggers on tables and views. See CREATE TRIGGER for more information.

Table 50-49. pg_trigger Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>tgrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table this trigger is on</td>
</tr>
<tr>
<td>tgname</td>
<td>name</td>
<td></td>
<td>Trigger name (must be unique among triggers of same table)</td>
</tr>
<tr>
<td>tgfoid</td>
<td>oid</td>
<td>pg_proc.oid</td>
<td>The function to be called</td>
</tr>
<tr>
<td>tgttype</td>
<td>int2</td>
<td></td>
<td>Bit mask identifying trigger firing conditions</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>----------------</td>
<td>------------</td>
<td>-------------</td>
</tr>
<tr>
<td>tgenabled</td>
<td>char</td>
<td></td>
<td>Controls in which session_replication_role modes the trigger fires. (O) = trigger fires in “origin” and “local” modes, (D) = trigger is disabled, (R) = trigger fires in “replica” mode, (A) = trigger fires always.</td>
</tr>
<tr>
<td>tgisinternal</td>
<td>bool</td>
<td></td>
<td>True if trigger is internally generated (usually, to enforce the constraint identified by tgconstraint)</td>
</tr>
<tr>
<td>tgconstrrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The table referenced by a referential integrity constraint</td>
</tr>
<tr>
<td>tgconstrindid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The index supporting a unique, primary key, referential integrity, or exclusion constraint</td>
</tr>
<tr>
<td>tgconstraint</td>
<td>oid</td>
<td>pg_constraint.oid</td>
<td>The pg_constraint entry associated with the trigger, if any</td>
</tr>
<tr>
<td>tgdeferrable</td>
<td>bool</td>
<td></td>
<td>True if constraint trigger is deferrable</td>
</tr>
<tr>
<td>tginitdeferred</td>
<td>bool</td>
<td></td>
<td>True if constraint trigger is initially deferred</td>
</tr>
<tr>
<td>tgnargs</td>
<td>int2</td>
<td></td>
<td>Number of argument strings passed to trigger function</td>
</tr>
<tr>
<td>tgattr</td>
<td>int2vector</td>
<td>pg_attribute.attnum</td>
<td>Column numbers, if trigger is column-specific; otherwise an empty array</td>
</tr>
<tr>
<td>tgargs</td>
<td>bytea</td>
<td></td>
<td>Argument strings to pass to trigger, each NULL-terminated</td>
</tr>
<tr>
<td>tgqual</td>
<td>pg_node_tree</td>
<td></td>
<td>Expression tree (in nodeToString() representation) for the trigger’s WHEN condition, or null if none</td>
</tr>
</tbody>
</table>
Currently, column-specific triggering is supported only for UPDATE events, and so \texttt{tgattr} is relevant only for that event type. \texttt{tgtype} might contain bits for other event types as well, but those are presumed to be table-wide regardless of what is in \texttt{tgattr}.

\textbf{Note:} When \texttt{tgconstraint} is nonzero, \texttt{tgconstrrelid}, \texttt{tgconstrindid}, \texttt{tdferrable}, and \texttt{tginitdeferred} are largely redundant with the referenced \texttt{pg_constraint} entry. However, it is possible for a non-deferrable trigger to be associated with a deferrable constraint: foreign key constraints can have some deferrable and some non-deferrable triggers.

\textbf{Note:} \texttt{pg_class.relhastriggers} must be true if a relation has any triggers in this catalog.

### 50.50. \texttt{pg\_ts\_config}

The \texttt{pg\_ts\_config} catalog contains entries representing text search configurations. A configuration specifies a particular text search parser and a list of dictionaries to use for each of the parser’s output token types. The parser is shown in the \texttt{pg\_ts\_config} entry, but the token-to-dictionary mapping is defined by subsidiary entries in \texttt{pg\_ts\_config\_map}.

PostgreSQL’s text search features are described at length in Chapter 12.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>\texttt{cfgname}</td>
<td>\texttt{name}</td>
<td></td>
<td>Text search configuration name</td>
</tr>
<tr>
<td>\texttt{cfgnamespace}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_namespace.oid}</td>
<td>The OID of the namespace that contains this configuration</td>
</tr>
<tr>
<td>\texttt{cfgowner}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_authid.oid}</td>
<td>Owner of the configuration</td>
</tr>
<tr>
<td>\texttt{cfgparser}</td>
<td>\texttt{oid}</td>
<td>\texttt{pg_ts_parser.oid}</td>
<td>The OID of the text search parser for this configuration</td>
</tr>
</tbody>
</table>

### 50.51. \texttt{pg\_ts\_config\_map}

The \texttt{pg\_ts\_config\_map} catalog contains entries showing which text search dictionaries should be consulted, and in what order, for each output token type of each text search configuration’s parser.

PostgreSQL’s text search features are described at length in Chapter 12.
Table 50-51. pg_ts_config_map Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>mapcfg</td>
<td>oid</td>
<td>pg_ts_config.oid</td>
<td>The OID of the pg_ts_config entry owning this map entry</td>
</tr>
<tr>
<td>maptokentype</td>
<td>integer</td>
<td></td>
<td>A token type emitted by the configuration's parser</td>
</tr>
<tr>
<td>mapseqno</td>
<td>integer</td>
<td></td>
<td>Order in which to consult this entry (lower mapseqnos first)</td>
</tr>
<tr>
<td>mapdict</td>
<td>oid</td>
<td>pg_ts_dict.oid</td>
<td>The OID of the text search dictionary to consult</td>
</tr>
</tbody>
</table>

50.52. pg_ts_dict

The pg_ts_dict catalog contains entries defining text search dictionaries. A dictionary depends on a text search template, which specifies all the implementation functions needed; the dictionary itself provides values for the user-settable parameters supported by the template. This division of labor allows dictionaries to be created by unprivileged users. The parameters are specified by a text string dictinitoption, whose format and meaning vary depending on the template.

PostgreSQL’s text search features are described at length in Chapter 12.

Table 50-52. pg_ts_dict Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>ddictname</td>
<td>name</td>
<td></td>
<td>Text search dictionary name</td>
</tr>
<tr>
<td>ddictnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this dictionary</td>
</tr>
<tr>
<td>ddictowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the dictionary</td>
</tr>
<tr>
<td>ddicttemplate</td>
<td>oid</td>
<td>pg_ts_template.oid</td>
<td>The OID of the text search template for this dictionary</td>
</tr>
<tr>
<td>ddictinitoption</td>
<td>text</td>
<td></td>
<td>Initialization option string for the template</td>
</tr>
</tbody>
</table>
50.53. **pg_ts_parser**

The `pg_ts_parser` catalog contains entries defining text search parsers. A parser is responsible for splitting input text into lexemes and assigning a token type to each lexeme. Since a parser must be implemented by C-language-level functions, creation of new parsers is restricted to database superusers.

PostgreSQL’s text search features are described at length in Chapter 12.

**Table 50-53. pg_ts_parser Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>prsname</td>
<td>name</td>
<td></td>
<td>Text search parser name</td>
</tr>
<tr>
<td>prsnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this parser</td>
</tr>
<tr>
<td>prsstart</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the parser’s startup function</td>
</tr>
<tr>
<td>prstoken</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the parser’s next-token function</td>
</tr>
<tr>
<td>prsend</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the parser’s shutdown function</td>
</tr>
<tr>
<td>prsheadline</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the parser’s headline function</td>
</tr>
<tr>
<td>prslextype</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>OID of the parser’s lexetype function</td>
</tr>
</tbody>
</table>

50.54. **pg_ts_template**

The `pg_ts_template` catalog contains entries defining text search templates. A template is the implementation skeleton for a class of text search dictionaries. Since a template must be implemented by C-language-level functions, creation of new templates is restricted to database superusers.

PostgreSQL’s text search features are described at length in Chapter 12.

**Table 50-54. pg_ts_template Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>tmplname</td>
<td>name</td>
<td></td>
<td>Text search template name</td>
</tr>
<tr>
<td>tmplnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this template</td>
</tr>
</tbody>
</table>
50.55. pg_type

The catalog `pg_type` stores information about data types. Base types and enum types (scalar types) are created with `CREATE TYPE`, and domains with `CREATE DOMAIN`. A composite type is automatically created for each table in the database, to represent the row structure of the table. It is also possible to create composite types with `CREATE TYPE AS`.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>typname</td>
<td>name</td>
<td></td>
<td>Data type name</td>
</tr>
<tr>
<td>typnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace that contains this type</td>
</tr>
<tr>
<td>typowner</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>Owner of the type</td>
</tr>
<tr>
<td>typlen</td>
<td>int2</td>
<td></td>
<td>For a fixed-size type, <code>typlen</code> is the number of bytes in the internal representation of the type. But for a variable-length type, <code>typlen</code> is negative. -1 indicates a “varlena” type (one that has a length word), -2 indicates a null-terminated C string.</td>
</tr>
</tbody>
</table>

Table 50-55. pg_type Columns
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>typbyval</td>
<td>bool</td>
<td></td>
<td>typbyval determines whether internal routines pass a value of this type by value or by reference. typbyval had better be false if typlen is not 1, 2, or 4 (or 8 on machines where Datum is 8 bytes). Variable-length types are always passed by reference. Note that typbyval can be false even if the length would allow pass-by-value.</td>
</tr>
<tr>
<td>typtype</td>
<td>char</td>
<td></td>
<td>typtype is b for a base type, c for a composite type (e.g., a table’s row type), d for a domain, e for an enum type, p for a pseudo-type, or r for a range type. See also typrelid and typbasetype.</td>
</tr>
<tr>
<td>typcategory</td>
<td>char</td>
<td></td>
<td>typcategory is an arbitrary classification of data types that is used by the parser to determine which implicit casts should be “preferred”. See also Table 50-56.</td>
</tr>
<tr>
<td>typispreferred</td>
<td>bool</td>
<td></td>
<td>True if the type is a preferred cast target within its typcategory</td>
</tr>
<tr>
<td>typisdefined</td>
<td>bool</td>
<td></td>
<td>True if the type is defined, false if this is a placeholder entry for a not-yet-defined type. When typisdefined is false, nothing except the type name, namespace, and OID can be relied on.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>---------</td>
<td>-------</td>
<td>------------------</td>
<td>---------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>typdelim</td>
<td>char</td>
<td></td>
<td>Character that separates two values of this type when parsing array input. Note that the delimiter is associated with the array element data type, not the array data type.</td>
</tr>
<tr>
<td>typrelid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>If this is a composite type (see typtype), then this column points to the pg_class entry that defines the corresponding table. (For a free-standing composite type, the pg_class entry doesn’t really represent a table, but it is needed anyway for the type’s pg_attribute entries to link to.) Zero for non-composite types.</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>typelem</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>If typelem is not 0, then it identifies another row in pg_type. The current type can then be subscripted like an array yielding values of type typelem. A “true” array type is variable length (typlen = -1), but some fixed-length (typlen &gt; 0) types also have nonzero typelem, for example name and point. If a fixed-length type has a typelem, then its internal representation must be some number of values of the typelem data type with no other data. Variable-length array types have a header defined by the array subroutines.</td>
</tr>
<tr>
<td>typarray</td>
<td>oid</td>
<td>pg_type.oid</td>
<td>If typarray is not 0, then it identifies another row in pg_type, which is the “true” array type having this type as element</td>
</tr>
<tr>
<td>typinput</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Input conversion function (text format)</td>
</tr>
<tr>
<td>typoutput</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Output conversion function (text format)</td>
</tr>
<tr>
<td>typreceive</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Input conversion function (binary format), or 0 if none</td>
</tr>
<tr>
<td>typsend</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Output conversion function (binary format), or 0 if none</td>
</tr>
<tr>
<td>typmodin</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Type modifier input function, or 0 if type does not support modifiers</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>typmodout</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Type modifier output function, or 0 to use the standard format</td>
</tr>
<tr>
<td>typanalyze</td>
<td>regproc</td>
<td>pg_proc.oid</td>
<td>Custom \texttt{ANALYZE} function, or 0 to use the standard function</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-------</td>
<td>------</td>
<td>------------</td>
<td>-------------</td>
</tr>
</tbody>
</table>
| typalign | char |            | typalign is the alignment required when storing a value of this type. It applies to storage on disk as well as most representations of the value inside PostgreSQL. When multiple values are stored consecutively, such as in the representation of a complete row on disk, padding is inserted before a datum of this type so that it begins on the specified boundary. The alignment reference is the beginning of the first datum in the sequence.

Possible values are:
- *c* = char alignment, i.e., no alignment needed.
- *s* = short alignment (2 bytes on most machines).
- *i* = int alignment (4 bytes on most machines).
- *d* = double alignment (8 bytes on many machines, but by no means all).

**Note:** For types used in system tables, it is critical that the size and alignment defined in `pg_type` agree with the way that the compiler will lay out the column in a structure representing a table row.

<p>| 2016 |</p>
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
</table>
| typstorage  | char  |            | typstorage tells for varlena types (those with \(\text{typlen} = -1\)) if the type is prepared for toasting and what the default strategy for attributes of this type should be. Possible values are:  
  - p: Value must always be stored plain.
  - e: Value can be stored in a “secondary” relation (if relation has one, see \(\text{pg\_class\.reltoastrelid}\)).
  - m: Value can be stored compressed inline.
  - x: Value can be stored compressed inline or stored in “secondary” storage.  
Note that m columns can also be moved out to secondary storage, but only as a last resort (e and x columns are moved first). |
<p>| typnotnull  |bool   |            | typnotnull represents a not-null constraint on a type. Used for domains only.                                                                 |
| typbasetype |oid    | pg_type.oid| If this is a domain (see typtype), then typbasetype identifies the type that this one is based on. Zero if this type is not a domain.             |</p>
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>typtypmod</td>
<td>int4</td>
<td></td>
<td>Domains use typtypmod to record the typmod to be applied to their base type (-1 if base type does not use a typmod). -1 if this type is not a domain.</td>
</tr>
<tr>
<td>typndims</td>
<td>int4</td>
<td></td>
<td>typndims is the number of array dimensions for a domain over an array (that is, typbasetype is an array type). Zero for types other than domains over array types.</td>
</tr>
<tr>
<td>typcollation</td>
<td>oid</td>
<td>pg_collation.oid</td>
<td>typcollation specifies the collation of the type. If the type does not support collations, this will be zero. A base type that supports collations will have DEFAULT_COLLATION_OID here. A domain over a collatable type can have some other collation OID, if one was specified for the domain.</td>
</tr>
<tr>
<td>typedefaultbin</td>
<td>pg_node_tree</td>
<td></td>
<td>If typedefaultbin is not null, it is the nodeToString() representation of a default expression for the type. This is only used for domains.</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>typdefault</td>
<td>text</td>
<td></td>
<td>typdefault is null if the type has no associated default value. If typdefaultbin is not null, typdefault must contain a human-readable version of the default expression represented by typdefaultbin. If typdefaultbin is null and typdefault is not, then typdefault is the external representation of the type’s default value, which can be fed to the type’s input converter to produce a constant.</td>
</tr>
<tr>
<td>typacl</td>
<td>aclitem[]</td>
<td></td>
<td>Access privileges; see GRANT and REVOKE for details.</td>
</tr>
</tbody>
</table>

Table 50-56 lists the system-defined values of typcategory. Any future additions to this list will also be upper-case ASCII letters. All other ASCII characters are reserved for user-defined categories.

### Table 50-56. typcategory Codes

<table>
<thead>
<tr>
<th>Code</th>
<th>Category</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>Array types</td>
</tr>
<tr>
<td>B</td>
<td>Boolean types</td>
</tr>
<tr>
<td>C</td>
<td>Composite types</td>
</tr>
<tr>
<td>D</td>
<td>Date/time types</td>
</tr>
<tr>
<td>E</td>
<td>Enum types</td>
</tr>
<tr>
<td>G</td>
<td>Geometric types</td>
</tr>
<tr>
<td>T</td>
<td>Network address types</td>
</tr>
<tr>
<td>N</td>
<td>Numeric types</td>
</tr>
<tr>
<td>P</td>
<td>Pseudo-types</td>
</tr>
<tr>
<td>R</td>
<td>Range types</td>
</tr>
<tr>
<td>S</td>
<td>String types</td>
</tr>
<tr>
<td>T</td>
<td>Timespan types</td>
</tr>
<tr>
<td>U</td>
<td>User-defined types</td>
</tr>
<tr>
<td>V</td>
<td>Bit-string types</td>
</tr>
<tr>
<td>X</td>
<td>unknown type</td>
</tr>
</tbody>
</table>

2019
50.56. pg_user_mapping

The catalog pg_user_mapping stores the mappings from local user to remote. Access to this catalog is restricted from normal users, use the view pg_user_mappings instead.

Table 50-57. pg_user_mapping Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>oid</td>
<td>oid</td>
<td></td>
<td>Row identifier (hidden attribute; must be explicitly selected)</td>
</tr>
<tr>
<td>umuser</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>OID of the local role being mapped, 0 if the user mapping is public</td>
</tr>
<tr>
<td>umserver</td>
<td>oid</td>
<td>pg_foreign_server.oid</td>
<td>The OID of the foreign server that contains this mapping</td>
</tr>
<tr>
<td>umoptions</td>
<td>text[]</td>
<td></td>
<td>User mapping specific options, as “keyword=value” strings</td>
</tr>
</tbody>
</table>

50.57. System Views

In addition to the system catalogs, PostgreSQL provides a number of built-in views. Some system views provide convenient access to some commonly used queries on the system catalogs. Other views provide access to internal server state.

The information schema (Chapter 35) provides an alternative set of views which overlap the functionality of the system views. Since the information schema is SQL-standard whereas the views described here are PostgreSQL-specific, it’s usually better to use the information schema if it provides all the information you need.

Table 50-58 lists the system views described here. More detailed documentation of each view follows below. There are some additional views that provide access to the results of the statistics collector; they are described in Table 28-2.

Except where noted, all the views described here are read-only.

Table 50-58. System Views

<table>
<thead>
<tr>
<th>View Name</th>
<th>Purpose</th>
</tr>
</thead>
<tbody>
<tr>
<td>pg_available_extensions</td>
<td>available extensions</td>
</tr>
<tr>
<td>pg_available_extension_versions</td>
<td>available versions of extensions</td>
</tr>
<tr>
<td>pg_config</td>
<td>compile-time configuration parameters</td>
</tr>
<tr>
<td>pg_cursors</td>
<td>open cursors</td>
</tr>
<tr>
<td>pg_file_settings</td>
<td>summary of configuration file contents</td>
</tr>
<tr>
<td>pg_group</td>
<td>groups of database users</td>
</tr>
<tr>
<td>pg_indexes</td>
<td>indexes</td>
</tr>
<tr>
<td>View Name</td>
<td>Purpose</td>
</tr>
<tr>
<td>---------------------------------</td>
<td>-------------------------------------------------------------------------</td>
</tr>
<tr>
<td>pg_locks</td>
<td>locks currently held or awaited</td>
</tr>
<tr>
<td>pg_matviews</td>
<td>materialized views</td>
</tr>
<tr>
<td>pg_policies</td>
<td>policies</td>
</tr>
<tr>
<td>pg_prepared_statements</td>
<td>prepared statements</td>
</tr>
<tr>
<td>pg_prepared_xacts</td>
<td>prepared transactions</td>
</tr>
<tr>
<td>pg_replication_origin_status</td>
<td>information about replication origins, including replication progress</td>
</tr>
<tr>
<td>pg_replication_slots</td>
<td>replication slot information</td>
</tr>
<tr>
<td>pg_roles</td>
<td>database roles</td>
</tr>
<tr>
<td>pg_rules</td>
<td>rules</td>
</tr>
<tr>
<td>pg_seclabels</td>
<td>security roles</td>
</tr>
<tr>
<td>pg_settings</td>
<td>parameter settings</td>
</tr>
<tr>
<td>pg_shadow</td>
<td>database users</td>
</tr>
<tr>
<td>pg_stats</td>
<td>planner statistics</td>
</tr>
<tr>
<td>pg_tables</td>
<td>tables</td>
</tr>
<tr>
<td>pg_timezone_abrevs</td>
<td>time zone abbreviations</td>
</tr>
<tr>
<td>pg_timezone_names</td>
<td>time zone names</td>
</tr>
<tr>
<td>pg_user</td>
<td>database users</td>
</tr>
<tr>
<td>pg_user_mappings</td>
<td>user mappings</td>
</tr>
<tr>
<td>pg_views</td>
<td>views</td>
</tr>
</tbody>
</table>

### 50.58. pg_available_extensions

The `pg_available_extensions` view lists the extensions that are available for installation. See also the `pg_extension` catalog, which shows the extensions currently installed.

**Table 50-59. pg_available_extensions Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>name</td>
<td>Extension name</td>
</tr>
<tr>
<td>default_version</td>
<td>text</td>
<td>Name of default version, or NULL if none is specified</td>
</tr>
<tr>
<td>installed_version</td>
<td>text</td>
<td>Currently installed version of the extension, or NULL if not installed</td>
</tr>
<tr>
<td>comment</td>
<td>text</td>
<td>Comment string from the extension’s control file</td>
</tr>
</tbody>
</table>

The `pg_available_extensions` view is read only.
50.59. **pg_available_extension_versions**

The `pg_available_extension_versions` view lists the specific extension versions that are available for installation. See also the `pg_extension` catalog, which shows the extensions currently installed.

**Table 50-60. pg_available_extension_versions Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>name</td>
<td>Extension name</td>
</tr>
<tr>
<td>version</td>
<td>text</td>
<td>Version name</td>
</tr>
<tr>
<td>installed</td>
<td>bool</td>
<td>True if this version of this extension is currently installed</td>
</tr>
<tr>
<td>superuser</td>
<td>bool</td>
<td>True if only superusers are allowed to install this extension</td>
</tr>
<tr>
<td>relocatable</td>
<td>bool</td>
<td>True if extension can be relocated to another schema</td>
</tr>
<tr>
<td>schema</td>
<td>name</td>
<td>Name of the schema that the extension must be installed into, or NULL if partially or fully relocatable</td>
</tr>
<tr>
<td>requires</td>
<td>name[]</td>
<td>Names of prerequisite extensions, or NULL if none</td>
</tr>
<tr>
<td>comment</td>
<td>text</td>
<td>Comment string from the extension’s control file</td>
</tr>
</tbody>
</table>

The `pg_available_extension_versions` view is read only.

50.60. **pg_config**

The view `pg_config` describes the compile-time configuration parameters of the currently installed version of PostgreSQL. It is intended, for example, to be used by software packages that want to interface to PostgreSQL to facilitate finding the required header files and libraries. It provides the same basic information as the `pg_config` PostgreSQL client application.

By default, the `pg_config` view can be read only by superusers.

**Table 50-61. pg_config Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>The parameter name</td>
</tr>
<tr>
<td>setting</td>
<td>text</td>
<td>The parameter value</td>
</tr>
</tbody>
</table>

50.61. **pg_cursors**

The `pg_cursors` view lists the cursors that are currently available. Cursors can be defined in several
ways:

- via the DECLARE statement in SQL
- via the Bind message in the frontend/backend protocol, as described in Section 51.2.3
- via the Server Programming Interface (SPI), as described in Section 45.1

The \texttt{pg\_cursors} view displays cursors created by any of these means. Cursors only exist for the duration of the transaction that defines them, unless they have been declared \texttt{WITH HOLD}. Therefore, non-holdable cursors are only present in the view until the end of their creating transaction.

\textbf{Note:} Cursors are used internally to implement some of the components of PostgreSQL, such as procedural languages. Therefore, the \texttt{pg\_cursors} view might include cursors that have not been explicitly created by the user.

### Table 50-62. \texttt{pg\_cursors} Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>The name of the cursor</td>
</tr>
<tr>
<td>statement</td>
<td>text</td>
<td>The verbatim query string submitted to declare this cursor</td>
</tr>
<tr>
<td>is_holdable</td>
<td>boolean</td>
<td>\texttt{true} if the cursor is holdable (that is, it can be accessed after the transaction that declared the cursor has committed); \texttt{false} otherwise</td>
</tr>
<tr>
<td>is_binary</td>
<td>boolean</td>
<td>\texttt{true} if the cursor was declared \texttt{BINARY}; \texttt{false} otherwise</td>
</tr>
<tr>
<td>is_scrollable</td>
<td>boolean</td>
<td>\texttt{true} if the cursor is scrollable (that is, it allows rows to be retrieved in a nonsequential manner); \texttt{false} otherwise</td>
</tr>
<tr>
<td>creation_time</td>
<td>timestamptz</td>
<td>The time at which the cursor was declared</td>
</tr>
</tbody>
</table>

The \texttt{pg\_cursors} view is read only.

### 50.62. \texttt{pg\_file\_settings}

The view \texttt{pg\_file\_settings} provides a summary of the contents of the server's configuration file(s). A row appears in this view for each “name = value” entry appearing in the files, with annotations indicating whether the value could be applied successfully. Additional row(s) may appear for problems not linked to a “name = value” entry, such as syntax errors in the files.

This view is helpful for checking whether planned changes in the configuration files will work, or for diagnosing a previous failure. Note that this view reports on the current contents of the files, not on what was last applied by the server. (The \texttt{pg\_settings} view is usually sufficient to determine that.)
By default, the `pg_file_settings` view can be read only by superusers.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>sourcefile</td>
<td>text</td>
<td>Full path name of the configuration file</td>
</tr>
<tr>
<td>sourceline</td>
<td>integer</td>
<td>Line number within the configuration file where the entry appears</td>
</tr>
<tr>
<td>seqno</td>
<td>integer</td>
<td>Order in which the entries are processed (1..n)</td>
</tr>
<tr>
<td>name</td>
<td>text</td>
<td>Configuration parameter name</td>
</tr>
<tr>
<td>setting</td>
<td>text</td>
<td>Value to be assigned to the parameter</td>
</tr>
<tr>
<td>applied</td>
<td>boolean</td>
<td>True if the value can be applied successfully</td>
</tr>
<tr>
<td>error</td>
<td>text</td>
<td>If not null, an error message indicating why this entry could not be applied</td>
</tr>
</tbody>
</table>

If the configuration file contains syntax errors or invalid parameter names, the server will not attempt to apply any settings from it, and therefore all the `applied` fields will read as false. In such a case there will be one or more rows with non-null `error` fields indicating the problem(s). Otherwise, individual settings will be applied if possible. If an individual setting cannot be applied (e.g., invalid value, or the setting cannot be changed after server start) it will have an appropriate message in the `error` field. Another way that an entry might have `applied = false` is that it is overridden by a later entry for the same parameter name; this case is not considered an error so nothing appears in the `error` field.

See Section 19.1 for more information about the various ways to change run-time parameters.

### 50.63. `pg_group`

The view `pg_group` exists for backwards compatibility: it emulates a catalog that existed in PostgreSQL before version 8.1. It shows the names and members of all roles that are marked as not `rolcanlogin`, which is an approximation to the set of roles that are being used as groups.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>groname</td>
<td>name</td>
<td><code>pg_authid.rolname</code></td>
<td>Name of the group</td>
</tr>
<tr>
<td>grosysid</td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>ID of this group</td>
</tr>
<tr>
<td>grolist</td>
<td>oid[]</td>
<td><code>pg_authid.oid</code></td>
<td>An array containing the IDs of the roles in this group</td>
</tr>
</tbody>
</table>
50.64. **pg_indexes**

The view `pg_indexes` provides access to useful information about each index in the database.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing table and index</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table the index is for</td>
</tr>
<tr>
<td>indexname</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of index</td>
</tr>
<tr>
<td>tablespace</td>
<td>name</td>
<td>pg_tablespace.spcname</td>
<td>Name of tablespace containing index (null if default for database)</td>
</tr>
<tr>
<td>indexdef</td>
<td>text</td>
<td></td>
<td>Index definition (a reconstructed CREATE INDEX command)</td>
</tr>
</tbody>
</table>

50.65. **pg_locks**

The view `pg_locks` provides access to information about the locks held by active processes within the database server. See Chapter 13 for more discussion of locking.

`pg_locks` contains one row per active lockable object, requested lock mode, and relevant process. Thus, the same lockable object might appear many times, if multiple processes are holding or waiting for locks on it. However, an object that currently has no locks on it will not appear at all.

There are several distinct types of lockable objects: whole relations (e.g., tables), individual pages of relations, individual tuples of relations, transaction IDs (both virtual and permanent IDs), and general database objects (identified by class OID and object OID, in the same way as in `pg_description` or `pg_depend`). Also, the right to extend a relation is represented as a separate lockable object. Also, “advisory” locks can be taken on numbers that have user-defined meanings.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>locktype</td>
<td>text</td>
<td></td>
<td>Type of the lockable object: relation, extend, page, tuple, transactionid, virtualxid, object, userlock, or advisory</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>-----------</td>
<td>----------</td>
<td>-----------------</td>
<td>----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>database</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>OID of the database in which the lock target exists, or zero if the target is a shared object, or null if the target is a transaction ID</td>
</tr>
<tr>
<td>relation</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>OID of the relation targeted by the lock, or null if the target is not a relation or part of a relation</td>
</tr>
<tr>
<td>page</td>
<td>integer</td>
<td></td>
<td>Page number targeted by the lock within the relation, or null if the target is not a relation page or tuple</td>
</tr>
<tr>
<td>tuple</td>
<td>smallint</td>
<td></td>
<td>Tuple number targeted by the lock within the page, or null if the target is not a tuple</td>
</tr>
<tr>
<td>virtualxid</td>
<td>text</td>
<td></td>
<td>Virtual ID of the transaction targeted by the lock, or null if the target is not a virtual transaction ID</td>
</tr>
<tr>
<td>transactionid</td>
<td>xid</td>
<td></td>
<td>ID of the transaction targeted by the lock, or null if the target is not a transaction ID</td>
</tr>
<tr>
<td>classid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>OID of the system catalog containing the lock target, or null if the target is not a general database object</td>
</tr>
<tr>
<td>objid</td>
<td>oid</td>
<td>any OID column</td>
<td>OID of the lock target within its system catalog, or null if the target is not a general database object</td>
</tr>
</tbody>
</table>
### Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>objsubid</td>
<td>smallint</td>
<td></td>
<td>Column number targeted by the lock (the classid and objid refer to the table itself), or zero if the target is some other general database object, or null if the target is not a general database object</td>
</tr>
<tr>
<td>virtualtransaction</td>
<td>text</td>
<td></td>
<td>Virtual ID of the transaction that is holding or awaiting this lock</td>
</tr>
<tr>
<td>pid</td>
<td>integer</td>
<td></td>
<td>Process ID of the server process holding or awaiting this lock, or null if the lock is held by a prepared transaction</td>
</tr>
<tr>
<td>mode</td>
<td>text</td>
<td></td>
<td>Name of the lock mode held or desired by this process (see Section 13.3.1 and Section 13.2.3)</td>
</tr>
<tr>
<td>granted</td>
<td>boolean</td>
<td></td>
<td>True if lock is held, false if lock is awaited</td>
</tr>
<tr>
<td>fastpath</td>
<td>boolean</td>
<td></td>
<td>True if lock was taken via fast path, false if taken via main lock table</td>
</tr>
</tbody>
</table>

*granted* is true in a row representing a lock held by the indicated process. False indicates that this process is currently waiting to acquire this lock, which implies that at least one other process is holding or waiting for a conflicting lock mode on the same lockable object. The waiting process will sleep until the other lock is released (or a deadlock situation is detected). A single process can be waiting to acquire at most one lock at a time.

Throughout running a transaction, a server process holds an exclusive lock on the transaction’s virtual transaction ID. If a permanent ID is assigned to the transaction (which normally happens only if the transaction changes the state of the database), it also holds an exclusive lock on the transaction’s permanent transaction ID until it ends. When a process finds it necessary to wait specifically for another transaction to end, it does so by attempting to acquire share lock on the other transaction’s ID (either virtual or permanent ID depending on the situation). That will succeed only when the other transaction terminates and releases its locks.

Although tuples are a lockable type of object, information about row-level locks is stored on disk, not in memory, and therefore row-level locks normally do not appear in this view. If a process is waiting for a row-level lock, it will usually appear in the view as waiting for the permanent transaction ID of the current holder of that row lock.
Advisory locks can be acquired on keys consisting of either a single bigint value or two integer values. A bigint key is displayed with its high-order half in the classid column, its low-order half in the objid column, and objsubid equal to 1. The original bigint value can be reassembled with the expression \((\text{classid}::\text{bigint} \ll 32) \mid \text{objid}::\text{bigint}\). Integer keys are displayed with the first key in the classid column, the second key in the objid column, and objsubid equal to 2. The actual meaning of the keys is up to the user. Advisory locks are local to each database, so the database column is meaningful for an advisory lock.

\textit{pg_locks} provides a global view of all locks in the database cluster, not only those relevant to the current database. Although its relation column can be joined against \textit{pg_class}.oid to identify locked relations, this will only work correctly for relations in the current database (those for which the database column is either the current database’s OID or zero).

The pid column can be joined to the pid column of the \textit{pg_stat_activity} view to get more information on the session holding or awaiting each lock, for example:

\begin{verbatim}
SELECT * FROM pg_locks pl LEFT JOIN pg_stat_activity psa
   ON pl.pid = psa.pid;
\end{verbatim}

Also, if you are using prepared transactions, the virtualtransaction column can be joined to the transaction column of the \textit{pg_prepared_xacts} view to get more information on prepared transactions that hold locks. (A prepared transaction can never be waiting for a lock, but it continues to hold the locks it acquired while running.) For example:

\begin{verbatim}
SELECT * FROM pg_locks pl LEFT JOIN pg_prepared_xacts ppx
   ON pl.virtualtransaction = '-1/' || ppx.transaction;
\end{verbatim}

While it is possible to obtain information about which processes block which other processes by joining \textit{pg_locks} against itself, this is very difficult to get right in detail. Such a query would have to encode knowledge about which lock modes conflict with which others. Worse, the \textit{pg_locks} view does not expose information about which processes are ahead of which others in lock wait queues, nor information about which processes are parallel workers running on behalf of which other client sessions. It is better to use the \textit{pg_blocking_pids()} function (see Table 9-59) to identify which process(es) a waiting process is blocked behind.

The \textit{pg_locks} view displays data from both the regular lock manager and the predicate lock manager, which are separate systems; in addition, the regular lock manager subdivides its locks into regular and fast-path locks. This data is not guaranteed to be entirely consistent. When the view is queried, data on fast-path locks (with \textit{fastpath} = \textit{true}) is gathered from each backend one at a time, without freezing the state of the entire lock manager, so it is possible for locks to be taken or released while information is gathered. Note, however, that these locks are known not to conflict with any other lock currently in place. After all backends have been queried for fast-path locks, the remainder of the regular lock manager is locked as a unit, and a consistent snapshot of all remaining locks is collected as an atomic action. After unlocking the regular lock manager, the predicate lock manager is similarly locked and all predicate locks are collected as an atomic action. Thus, with the exception of fast-path locks, each lock manager will deliver a consistent set of results, but as we do not lock both lock managers simultaneously, it is possible for locks to be taken or released after we interrogate the regular lock manager and before we interrogate the predicate lock manager.

Locking the regular and/or predicate lock manager could have some impact on database performance if this view is very frequently accessed. The locks are held only for the minimum amount of time necessary to obtain data from the lock managers, but this does not completely eliminate the possibility of a performance impact.
50.66. `pg_matviews`  

The view `pg_matviews` provides access to useful information about each materialized view in the database.

**Table 50-67. `pg_matviews` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td><code>pg_namespace.nspname</code></td>
<td>Name of schema containing materialized view</td>
</tr>
<tr>
<td>matviewname</td>
<td>name</td>
<td><code>pg_class.relname</code></td>
<td>Name of materialized view</td>
</tr>
<tr>
<td>matviewowner</td>
<td>name</td>
<td><code>pg_authid.rolname</code></td>
<td>Name of materialized view’s owner</td>
</tr>
<tr>
<td>tablespace</td>
<td>name</td>
<td><code>pg_tablespace.spcname</code></td>
<td>Name of tablespace containing materialized view (null if default for database)</td>
</tr>
<tr>
<td>hasindexes</td>
<td>boolean</td>
<td></td>
<td>True if materialized view has (or recently had) any indexes</td>
</tr>
<tr>
<td>ispopulated</td>
<td>boolean</td>
<td></td>
<td>True if materialized view is currently populated</td>
</tr>
<tr>
<td>definition</td>
<td>text</td>
<td></td>
<td>Materialized view definition (a reconstructed SELECT query)</td>
</tr>
</tbody>
</table>

50.67. `pg_policies`  

The view `pg_policies` provides access to useful information about each row-level security policy in the database.

**Table 50-68. `pg_policies` Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td><code>pg_namespace.nspname</code></td>
<td>Name of schema containing table policy is on</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td><code>pg_class.relname</code></td>
<td>Name of table policy is on</td>
</tr>
<tr>
<td>policyname</td>
<td>name</td>
<td><code>pg_policy.polname</code></td>
<td>Name of policy</td>
</tr>
<tr>
<td>roles</td>
<td>name[]</td>
<td></td>
<td>The roles to which this policy applies</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>cmd</td>
<td>text</td>
<td></td>
<td>The command type to which the policy is applied</td>
</tr>
<tr>
<td>qual</td>
<td>text</td>
<td></td>
<td>The expression added to the security barrier qualifications for queries that this policy applies to</td>
</tr>
<tr>
<td>with_check</td>
<td>text</td>
<td></td>
<td>The expression added to the WITH CHECK qualifications for queries that attempt to add rows to this table</td>
</tr>
</tbody>
</table>

50.68. `pg_prepared_statements`

The `pg_prepared_statements` view displays all the prepared statements that are available in the current session. See PREPARE for more information about prepared statements.

`pg_prepared_statements` contains one row for each prepared statement. Rows are added to the view when a new prepared statement is created and removed when a prepared statement is released (for example, via the DEALLOCATE command).

Table 50-69. `pg_prepared_statements` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>The identifier of the prepared statement</td>
</tr>
<tr>
<td>statement</td>
<td>text</td>
<td>The query string submitted by the client to create this prepared statement. For prepared statements created via SQL, this is the PREPARE statement submitted by the client. For prepared statements created via the frontend/backend protocol, this is the text of the prepared statement itself.</td>
</tr>
<tr>
<td>prepare_time</td>
<td>timestamptz</td>
<td>The time at which the prepared statement was created</td>
</tr>
</tbody>
</table>
### 50.69. pg_prepared_xacts

The view `pg_prepared_xacts` displays information about transactions that are currently prepared for two-phase commit (see `PREPARE TRANSACTION` for details).

`pg_prepared_xacts` contains one row per prepared transaction. An entry is removed when the transaction is committed or rolled back.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>transaction</td>
<td>xid</td>
<td>Numeric transaction identifier of the prepared transaction</td>
</tr>
<tr>
<td>gid</td>
<td>text</td>
<td>Global transaction identifier that was assigned to the transaction</td>
</tr>
<tr>
<td>prepared</td>
<td>timestamp with time zone</td>
<td>Time at which the transaction was prepared for commit</td>
</tr>
<tr>
<td>owner</td>
<td>name</td>
<td>Name of the user that executed the transaction</td>
</tr>
<tr>
<td>database</td>
<td>name</td>
<td>Name of the database in which the transaction was executed</td>
</tr>
</tbody>
</table>

When the `pg_prepared_xacts` view is accessed, the internal transaction manager data structures are momentarily locked, and a copy is made for the view to display. This ensures that the view produces a consistent set of results, while not blocking normal operations longer than necessary. Nonetheless
there could be some impact on database performance if this view is frequently accessed.

50.70. \texttt{pg\_replication\_origin\_status}

The \texttt{pg\_replication\_origin\_status} view contains information about how far replay for a certain origin has progressed. For more on replication origins see Chapter 48.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>local_id</td>
<td>Oid</td>
<td>\texttt{pg_replication_origin_status}</td>
<td>internal node identifier</td>
</tr>
<tr>
<td>external_id</td>
<td>text</td>
<td>\texttt{pg_replication_origin_status}</td>
<td>external node identifier</td>
</tr>
<tr>
<td>remote_lsn</td>
<td>\texttt{pg_lsn}</td>
<td></td>
<td>The origin node’s LSN up to which data has been replicated.</td>
</tr>
<tr>
<td>local_lsn</td>
<td>\texttt{pg_lsn}</td>
<td></td>
<td>This node’s LSN at which \texttt{remote_lsn} has been replicated. This node’s LSN at which \texttt{remote_lsn} has been replicated. Used to flush commit records before persisting data to disk when using asynchronous commits.</td>
</tr>
</tbody>
</table>

50.71. \texttt{pg\_replication\_slots}

The \texttt{pg\_replication\_slots} view provides a listing of all replication slots that currently exist on the database cluster, along with their current state.

For more on replication slots, see Section 26.2.6 and Chapter 47.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>slot_name</td>
<td>name</td>
<td></td>
<td>A unique, cluster-wide identifier for the replication slot</td>
</tr>
<tr>
<td>plugin</td>
<td>name</td>
<td></td>
<td>The base name of the shared object containing the output plugin this logical slot is using, or null for physical slots.</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>---------</td>
<td>------------</td>
<td>---------------------------------------------------------</td>
</tr>
<tr>
<td>slot_type</td>
<td>text</td>
<td></td>
<td>The slot type - physical or logical</td>
</tr>
<tr>
<td>datoid</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>The OID of the database this slot is associated with, or null. Only logical slots have an associated database.</td>
</tr>
<tr>
<td>database</td>
<td>text</td>
<td>pg_database.datname</td>
<td>The name of the database this slot is associated with, or null. Only logical slots have an associated database.</td>
</tr>
<tr>
<td>active</td>
<td>boolean</td>
<td></td>
<td>True if this slot is currently actively being used</td>
</tr>
<tr>
<td>active_pid</td>
<td>integer</td>
<td></td>
<td>The process ID of the session using this slot if the slot is currently actively being used. NULL if inactive.</td>
</tr>
<tr>
<td>xmin</td>
<td>xid</td>
<td></td>
<td>The oldest transaction that this slot needs the database to retain. VACUUM cannot remove tuples deleted by any later transaction.</td>
</tr>
<tr>
<td>catalog_xmin</td>
<td>xid</td>
<td></td>
<td>The oldest transaction affecting the system catalogs that this slot needs the database to retain. VACUUM cannot remove catalog tuples deleted by any later transaction.</td>
</tr>
<tr>
<td>restart_lsn</td>
<td>pg_lsn</td>
<td></td>
<td>The address (LSN) of oldest WAL which still might be required by the consumer of this slot and thus won’t be automatically removed during checkpoints.</td>
</tr>
</tbody>
</table>
## 50.72. `pg_roles`

The view `pg_roles` provides access to information about database roles. This is simply a publicly readable view of `pg_authid` that blanks out the password field.

This view explicitly exposes the OID column of the underlying table, since that is needed to do joins to other catalogs.

### Table 50-73. `pg_roles` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rolname</td>
<td>name</td>
<td></td>
<td>Role name</td>
</tr>
<tr>
<td>rolsuper</td>
<td>bool</td>
<td></td>
<td>Role has superuser privileges</td>
</tr>
<tr>
<td>rolinherit</td>
<td>bool</td>
<td></td>
<td>Role automatically inherits privileges of roles it is a member of</td>
</tr>
<tr>
<td>rolcreaterole</td>
<td>bool</td>
<td></td>
<td>Role can create more roles</td>
</tr>
<tr>
<td>rolcreatedb</td>
<td>bool</td>
<td></td>
<td>Role can create databases</td>
</tr>
<tr>
<td>rolcanlogin</td>
<td>bool</td>
<td></td>
<td>Role can log in. That is, this role can be given as the initial session authorization identifier</td>
</tr>
<tr>
<td>rolreplication</td>
<td>bool</td>
<td></td>
<td>Role is a replication role. A replication role can initiate replication connections and create and drop replication slots.</td>
</tr>
<tr>
<td>rolconnlimit</td>
<td>int4</td>
<td></td>
<td>For roles that can log in, this sets maximum number of concurrent connections this role can make. -1 means no limit.</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>rolpassword</td>
<td>text</td>
<td></td>
<td>Not the password (always reads as ************)</td>
</tr>
<tr>
<td>rolvaliduntil</td>
<td>timestamptz</td>
<td></td>
<td>Password expiry time (only used for password authentication); null if no expiration</td>
</tr>
<tr>
<td>rolbypassrls</td>
<td>bool</td>
<td></td>
<td>Role bypasses every row level security policy, see Section 5.7 for more information.</td>
</tr>
<tr>
<td>rolconfig</td>
<td>text[]</td>
<td></td>
<td>Role-specific defaults for run-time configuration variables</td>
</tr>
<tr>
<td>oid</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>ID of role</td>
</tr>
</tbody>
</table>

50.73. pg_rules

The view pg_rules provides access to useful information about query rewrite rules.

Table 50-74. pg_rules Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing table</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table the rule is for</td>
</tr>
<tr>
<td>rulename</td>
<td>name</td>
<td>pg_rewrite.rulename</td>
<td>Name of rule</td>
</tr>
<tr>
<td>definition</td>
<td>text</td>
<td></td>
<td>Rule definition (a reconstructed creation command)</td>
</tr>
</tbody>
</table>

The pg_rules view excludes the ON SELECT rules of views and materialized views; those can be seen in pg_views and pg_matviews.

50.74. pg_seclabels

The view pg_seclabels provides information about security labels. It is an easier-to-query version of the pg_seclabel catalog.

Table 50-75. pg_seclabels Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>------------</td>
<td>---------</td>
<td>------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>objoid</td>
<td>oid</td>
<td>any OID column</td>
<td>The OID of the object this security label pertains to</td>
</tr>
<tr>
<td>classoid</td>
<td>oid</td>
<td>pg_class.oid</td>
<td>The OID of the system catalog this object appears in</td>
</tr>
<tr>
<td>objsubid</td>
<td>int4</td>
<td></td>
<td>For a security label on a table column, this is the column number (the objoid and classoid refer to the table itself). For all other object types, this column is zero.</td>
</tr>
<tr>
<td>objtype</td>
<td>text</td>
<td></td>
<td>The type of object to which this label applies, as text.</td>
</tr>
<tr>
<td>objnamespace</td>
<td>oid</td>
<td>pg_namespace.oid</td>
<td>The OID of the namespace for this object, if applicable; otherwise NULL.</td>
</tr>
<tr>
<td>objname</td>
<td>text</td>
<td></td>
<td>The name of the object to which this label applies, as text.</td>
</tr>
<tr>
<td>provider</td>
<td>text</td>
<td>pg_seclabel.provider</td>
<td>The label provider associated with this label.</td>
</tr>
<tr>
<td>label</td>
<td>text</td>
<td>pg_seclabel.label</td>
<td>The security label applied to this object.</td>
</tr>
</tbody>
</table>

### 50.75. pg_settings

The view `pg_settings` provides access to run-time parameters of the server. It is essentially an alternative interface to the SHOW and SET commands. It also provides access to some facts about each parameter that are not directly available from SHOW, such as minimum and maximum values.

#### Table 50-76. pg_settings Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>Run-time configuration parameter name</td>
</tr>
<tr>
<td>setting</td>
<td>text</td>
<td>Current value of the parameter</td>
</tr>
<tr>
<td>unit</td>
<td>text</td>
<td>Implicit unit of the parameter</td>
</tr>
<tr>
<td>category</td>
<td>text</td>
<td>Logical group of the parameter</td>
</tr>
<tr>
<td>short_desc</td>
<td>text</td>
<td>A brief description of the parameter</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>-------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>extra_desc</td>
<td>text</td>
<td>Additional, more detailed, description of the parameter</td>
</tr>
<tr>
<td>context</td>
<td>text</td>
<td>Context required to set the parameter’s value (see below)</td>
</tr>
<tr>
<td>vartype</td>
<td>text</td>
<td>Parameter type (bool, enum, integer, real, or string)</td>
</tr>
<tr>
<td>source</td>
<td>text</td>
<td>Source of the current parameter value</td>
</tr>
<tr>
<td>min_val</td>
<td>text</td>
<td>Minimum allowed value of the parameter (null for non-numeric values)</td>
</tr>
<tr>
<td>max_val</td>
<td>text</td>
<td>Maximum allowed value of the parameter (null for non-numeric values)</td>
</tr>
<tr>
<td>enumvals</td>
<td>text[]</td>
<td>Allowed values of an enum parameter (null for non-enum values)</td>
</tr>
<tr>
<td>boot_val</td>
<td>text</td>
<td>Parameter value assumed at server startup if the parameter is not otherwise set</td>
</tr>
<tr>
<td>reset_val</td>
<td>text</td>
<td>Value that RESET would reset the parameter to in the current session</td>
</tr>
<tr>
<td>sourcefile</td>
<td>text</td>
<td>Configuration file the current value was set in (null for values set from sources other than configuration files, or when examined by a non-superuser); helpful when using include directives in configuration files</td>
</tr>
<tr>
<td>sourceline</td>
<td>integer</td>
<td>Line number within the configuration file the current value was set at (null for values set from sources other than configuration files, or when examined by a non-superuser)</td>
</tr>
<tr>
<td>pending_restart</td>
<td>boolean</td>
<td>true if the value has been changed in the configuration file but needs a restart; or false otherwise.</td>
</tr>
</tbody>
</table>

There are several possible values of context. In order of decreasing difficulty of changing the setting, they are:

**internal**

These settings cannot be changed directly; they reflect internally determined values. Some of them may be adjustable by rebuilding the server with different configuration options, or by changing options supplied to initdb.
postmaster

These settings can only be applied when the server starts, so any change requires restarting the server. Values for these settings are typically stored in the `postgresql.conf` file, or passed on the command line when starting the server. Of course, settings with any of the lower context types can also be set at server start time.

sighup

Changes to these settings can be made in `postgresql.conf` without restarting the server. Send a SIGHUP signal to the postmaster to cause it to re-read `postgresql.conf` and apply the changes. The postmaster will also forward the SIGHUP signal to its child processes so that they all pick up the new value.

superuser-backend

Changes to these settings can be made in `postgresql.conf` without restarting the server. They can also be set for a particular session in the connection request packet (for example, via libpq’s `PGOPTIONS` environment variable), but only if the connecting user is a superuser. However, these settings never change in a session after it is started. If you change them in `postgresql.conf`, send a SIGHUP signal to the postmaster to cause it to re-read `postgresql.conf`. The new values will only affect subsequently-launched sessions.

backend

Changes to these settings can be made in `postgresql.conf` without restarting the server. They can also be set for a particular session in the connection request packet (for example, via libpq’s `PGOPTIONS` environment variable); any user can make such a change for their session. However, these settings never change in a session after it is started. If you change them in `postgresql.conf`, send a SIGHUP signal to the postmaster to cause it to re-read `postgresql.conf`. The new values will only affect subsequently-launched sessions.

superuser

These settings can be set from `postgresql.conf`, or within a session via the `SET` command; but only superusers can change them via `SET`. Changes in `postgresql.conf` will affect existing sessions only if no session-local value has been established with `SET`.

user

These settings can be set from `postgresql.conf`, or within a session via the `SET` command. Any user is allowed to change their session-local value. Changes in `postgresql.conf` will affect existing sessions only if no session-local value has been established with `SET`.

See Section 19.1 for more information about the various ways to change these parameters.

The `pg_settings` view cannot be inserted into or deleted from, but it can be updated. An `UPDATE` applied to a row of `pg_settings` is equivalent to executing the `SET` command on that named parameter. The change only affects the value used by the current session. If an `UPDATE` is issued within a transaction that is later aborted, the effects of the `UPDATE` command disappear when the transaction is rolled back. Once the surrounding transaction is committed, the effects will persist until the end of the session, unless overridden by another `UPDATE` or `SET`.

50.76. pg_shadow

The view `pg_shadow` exists for backwards compatibility: it emulates a catalog that existed in PostgreSQL before version 8.1. It shows properties of all roles that are marked as `rolcanlogin` in `pg_authid`. 
The name stems from the fact that this table should not be readable by the public since it contains passwords. `pg_user` is a publicly readable view on `pg_shadow` that blanks out the password field.

### Table 50-77. pg_shadow Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>usename</td>
<td>name</td>
<td><code>pg_authid.rolname</code></td>
<td>User name</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td><code>pg_authid.oid</code></td>
<td>ID of this user</td>
</tr>
<tr>
<td>usecreatedb</td>
<td>bool</td>
<td></td>
<td>User can create databases</td>
</tr>
<tr>
<td>usesuper</td>
<td>bool</td>
<td></td>
<td>User is a superuser</td>
</tr>
<tr>
<td>userepl</td>
<td>bool</td>
<td></td>
<td>User can initiate streaming replication and put the system in and out of backup mode.</td>
</tr>
<tr>
<td>usebypassrls</td>
<td>bool</td>
<td></td>
<td>User bypasses every row level security policy, see Section 5.7 for more information.</td>
</tr>
<tr>
<td>passwd</td>
<td>text</td>
<td></td>
<td>Password (possibly encrypted); null if none. See <code>pg_authid</code> for details of how encrypted passwords are stored.</td>
</tr>
<tr>
<td>valuntil</td>
<td>abstime</td>
<td></td>
<td>Password expiry time (only used for password authentication)</td>
</tr>
<tr>
<td>useconfig</td>
<td>text[]</td>
<td></td>
<td>Session defaults for run-time configuration variables</td>
</tr>
</tbody>
</table>

### 50.77. pg_stats

The view `pg_stats` provides access to the information stored in the `pg_statistic` catalog. This view allows access only to rows of `pg_statistic` that correspond to tables the user has permission to read, and therefore it is safe to allow public read access to this view.

`pg_stats` is also designed to present the information in a more readable format than the underlying catalog — at the cost that its schema must be extended whenever new slot types are defined for `pg_statistic`.

### Table 50-78. pg_stats Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>---------------</td>
<td>---------</td>
<td>--------------------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing table</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table</td>
</tr>
<tr>
<td>attname</td>
<td>name</td>
<td>pg_attribute.attname</td>
<td>Name of the column described by this row</td>
</tr>
<tr>
<td>inherited</td>
<td>bool</td>
<td></td>
<td>If true, this row includes inheritance child columns, not just the values in the specified table</td>
</tr>
<tr>
<td>null_frac</td>
<td>real</td>
<td></td>
<td>Fraction of column entries that are null</td>
</tr>
<tr>
<td>avg_width</td>
<td>integer</td>
<td></td>
<td>Average width in bytes of column’s entries</td>
</tr>
<tr>
<td>n_distinct</td>
<td>real</td>
<td></td>
<td>If greater than zero, the estimated number of distinct values in the column. If less than zero, the negative of the number of distinct values divided by the number of rows. (The negated form is used when \texttt{ANALYZE} believes that the number of distinct values is likely to increase as the table grows; the positive form is used when the column seems to have a fixed number of possible values.) For example, -1 indicates a unique column in which the number of distinct values is the same as the number of rows.</td>
</tr>
<tr>
<td>most_common_vals</td>
<td>anyarray</td>
<td></td>
<td>A list of the most common values in the column. (Null if no values seem to be more common than any others.)</td>
</tr>
</tbody>
</table>

2040
<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>most_common_freqs</td>
<td>real[]</td>
<td></td>
<td>A list of the frequencies of the most common values, i.e., number of occurrences of each divided by total number of rows. (Null when most_common_vals is.)</td>
</tr>
<tr>
<td>histogram_bounds</td>
<td>anyarray</td>
<td></td>
<td>A list of values that divide the column’s values into groups of approximately equal population. The values in most_common_vals, if present, are omitted from this histogram calculation. (This column is null if the column data type does not have a &lt; operator or if the most_common_vals list accounts for the entire population.)</td>
</tr>
<tr>
<td>correlation</td>
<td>real</td>
<td></td>
<td>Statistical correlation between physical row ordering and logical ordering of the column values. This ranges from -1 to +1. When the value is near -1 or +1, an index scan on the column will be estimated to be cheaper than when it is near zero, due to reduction of random access to the disk. (This column is null if the column data type does not have a &lt; operator.)</td>
</tr>
<tr>
<td>most_common_elems</td>
<td>anyarray</td>
<td></td>
<td>A list of non-null element values most often appearing within values of the column. (Null for scalar types.)</td>
</tr>
</tbody>
</table>
Chapter 50. System Catalogs

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>most_common_elem_freqs</td>
<td>real[]</td>
<td></td>
<td>A list of the frequencies of the most common element values, i.e., the fraction of rows containing at least one instance of the given value. Two or three additional values follow the per-element frequencies; these are the minimum and maximum of the preceding per-element frequencies, and optionally the frequency of null elements. (Null when most_commonelems is.)</td>
</tr>
<tr>
<td>elem_count_histogram</td>
<td>real[]</td>
<td></td>
<td>A histogram of the counts of distinct non-null element values within the values of the column, followed by the average number of distinct non-null elements. (Null for scalar types.)</td>
</tr>
</tbody>
</table>

The maximum number of entries in the array fields can be controlled on a column-by-column basis using the ALTER TABLE SET STATISTICS command, or globally by setting the default_statistics_target run-time parameter.

50.78. pg_tables

The view pg_tables provides access to useful information about each table in the database.

Table 50-79. pg_tables Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing table</td>
</tr>
<tr>
<td>tablename</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of table</td>
</tr>
<tr>
<td>tableowner</td>
<td>name</td>
<td>pg_authid.rolname</td>
<td>Name of table’s owner</td>
</tr>
<tr>
<td>tablespace</td>
<td>name</td>
<td>pg_tablespace.spcname</td>
<td>Name of tablespace containing table (null if default for database)</td>
</tr>
</tbody>
</table>
50.79. **pg_timezone_abrevs**

The view `pg_timezone_abrevs` provides a list of time zone abbreviations that are currently recognized by the datetime input routines. The contents of this view change when the timezone_abreviations run-time parameter is modified.

Table **50-80. pg_timezone_abrevs Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>abbrev</td>
<td>text</td>
<td>Time zone abbreviation</td>
</tr>
<tr>
<td>utc_offset</td>
<td>interval</td>
<td>Offset from UTC (positive means east of Greenwich)</td>
</tr>
<tr>
<td>is_dst</td>
<td>boolean</td>
<td>True if this is a daylight-savings abbreviation</td>
</tr>
</tbody>
</table>

While most timezone abbreviations represent fixed offsets from UTC, there are some that have historically varied in value (see Section B.4 for more information). In such cases this view presents their current meaning.

50.80. **pg_timezone_names**

The view `pg_timezone_names` provides a list of time zone names that are recognized by `SET TIMEZONE`, along with their associated abbreviations, UTC offsets, and daylight-savings status. (Technically, PostgreSQL does not use UTC because leap seconds are not handled.) Unlike the abbreviations shown in `pg_timezone_abrevs`, many of these names imply a set of daylight-savings transition date rules. Therefore, the associated information changes across local DST boundaries. The displayed information is computed based on the current value of `CURRENT_TIMESTAMP`.

Table **50-81. pg_timezone_names Columns**

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>text</td>
<td>Time zone name</td>
</tr>
<tr>
<td>abbrev</td>
<td>text</td>
<td>Time zone abbreviation</td>
</tr>
<tr>
<td>utc_offset</td>
<td>interval</td>
<td>Offset from UTC (positive means east of Greenwich)</td>
</tr>
</tbody>
</table>
### 50.81. pg_user

The view `pg_user` provides access to information about database users. This is simply a publicly readable view of `pg_shadow` that blanks out the password field.

Table 50-82. `pg_user` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>usename</td>
<td>name</td>
<td>User name</td>
</tr>
<tr>
<td>usesysid</td>
<td>oid</td>
<td>ID of this user</td>
</tr>
<tr>
<td>usecreatedb</td>
<td>bool</td>
<td>User can create databases</td>
</tr>
<tr>
<td>usesuper</td>
<td>bool</td>
<td>User is a superuser</td>
</tr>
<tr>
<td>userepl</td>
<td>bool</td>
<td>User can initiate streaming replication and put the system in and out of backup mode.</td>
</tr>
<tr>
<td>usebypassrls</td>
<td>bool</td>
<td>User bypasses every row level security policy, see Section 5.7 for more information.</td>
</tr>
<tr>
<td>passwd</td>
<td>text</td>
<td>Not the password (always reads as ********)</td>
</tr>
<tr>
<td>valuntil</td>
<td>abstime</td>
<td>Password expiry time (only used for password authentication)</td>
</tr>
<tr>
<td>useconfig</td>
<td>text[]</td>
<td>Session defaults for run-time configuration variables</td>
</tr>
</tbody>
</table>

### 50.82. pg_user_mappings

The view `pg_user_mappings` provides access to information about user mappings. This is essentially a publicly readable view of `pg_user_mapping` that leaves out the options field if the user has no rights to use it.

Table 50-83. `pg_user_mappings` Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>umid</td>
<td>oid</td>
<td><code>pg_user_mapping.oid</code></td>
<td>OID of the user mapping</td>
</tr>
<tr>
<td>srvid</td>
<td>oid</td>
<td><code>pg_foreign_server.oid</code></td>
<td>The OID of the foreign server that contains this mapping</td>
</tr>
</tbody>
</table>
50.83. pg_views

The view pg_views provides access to useful information about each view in the database.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>schemaname</td>
<td>name</td>
<td>pg_namespace.nspname</td>
<td>Name of schema containing view</td>
</tr>
<tr>
<td>viewname</td>
<td>name</td>
<td>pg_class.relname</td>
<td>Name of view</td>
</tr>
<tr>
<td>viewowner</td>
<td>name</td>
<td>pg_authid.rolname</td>
<td>Name of view’s owner</td>
</tr>
<tr>
<td>definition</td>
<td>text</td>
<td></td>
<td>View definition (a reconstructed SELECT query)</td>
</tr>
</tbody>
</table>
Chapter 51. Frontend/Backend Protocol

PostgreSQL uses a message-based protocol for communication between frontends and backends (clients and servers). The protocol is supported over TCP/IP and also over Unix-domain sockets. Port number 5432 has been registered with IANA as the customary TCP port number for servers supporting this protocol, but in practice any non-privileged port number can be used.

This document describes version 3.0 of the protocol, implemented in PostgreSQL 7.4 and later. For descriptions of the earlier protocol versions, see previous releases of the PostgreSQL documentation. A single server can support multiple protocol versions. The initial startup-request message tells the server which protocol version the client is attempting to use. If the major version requested by the client is not supported by the server, the connection will be rejected (for example, this would occur if the client requested protocol version 4.0, which does not exist as of this writing). If the minor version requested by the client is not supported by the server (e.g. the client requests version 3.1, but the server supports only 3.0), the server may either reject the connection or may respond with a NegotiateProtocolVersion message containing the highest minor protocol version which it supports. The client may then choose either to continue with the connection using the specified protocol version or to abort the connection.

In order to serve multiple clients efficiently, the server launches a new “backend” process for each client. In the current implementation, a new child process is created immediately after an incoming connection is detected. This is transparent to the protocol, however. For purposes of the protocol, the terms “backend” and “server” are interchangeable; likewise “frontend” and “client” are interchangeable.

51.1. Overview

The protocol has separate phases for startup and normal operation. In the startup phase, the frontend opens a connection to the server and authenticates itself to the satisfaction of the server. (This might involve a single message, or multiple messages depending on the authentication method being used.) If all goes well, the server then sends status information to the frontend, and finally enters normal operation. Except for the initial startup-request message, this part of the protocol is driven by the server.

During normal operation, the frontend sends queries and other commands to the backend, and the backend sends back query results and other responses. There are a few cases (such as NOTIFY) wherein the backend will send unsolicited messages, but for the most part this portion of a session is driven by frontend requests.

Termination of the session is normally by frontend choice, but can be forced by the backend in certain cases. In any case, when the backend closes the connection, it will roll back any open (incomplete) transaction before exiting.

Within normal operation, SQL commands can be executed through either of two sub-protocols. In the “simple query” protocol, the frontend just sends a textual query string, which is parsed and immediately executed by the backend. In the “extended query” protocol, processing of queries is separated into multiple steps: parsing, binding of parameter values, and execution. This offers flexibility and performance benefits, at the cost of extra complexity.

Normal operation has additional sub-protocols for special operations such as COPY.
51.1.1. Messaging Overview

All communication is through a stream of messages. The first byte of a message identifies the message type, and the next four bytes specify the length of the rest of the message (this length count includes itself, but not the message-type byte). The remaining contents of the message are determined by the message type. For historical reasons, the very first message sent by the client (the startup message) has no initial message-type byte.

To avoid losing synchronization with the message stream, both servers and clients typically read an entire message into a buffer (using the byte count) before attempting to process its contents. This allows easy recovery if an error is detected while processing the contents. In extreme situations (such as not having enough memory to buffer the message), the receiver can use the byte count to determine how much input to skip before it resumes reading messages.

Conversely, both servers and clients must take care never to send an incomplete message. This is commonly done by marshaling the entire message in a buffer before beginning to send it. If a communications failure occurs partway through sending or receiving a message, the only sensible response is to abandon the connection, since there is little hope of recovering message-boundary synchronization.

51.1.2. Extended Query Overview

In the extended-query protocol, execution of SQL commands is divided into multiple steps. The state retained between steps is represented by two types of objects: prepared statements and portals. A prepared statement represents the result of parsing and semantic analysis of a textual query string. A prepared statement is not in itself ready to execute, because it might lack specific values for parameters. A portal represents a ready-to-execute or already-partially-executed statement, with any missing parameter values filled in. (For SELECT statements, a portal is equivalent to an open cursor, but we choose to use a different term since cursors don’t handle non-SELECT statements.)

The overall execution cycle consists of a parse step, which creates a prepared statement from a textual query string; a bind step, which creates a portal given a prepared statement and values for any needed parameters; and an execute step that runs a portal’s query. In the case of a query that returns rows (SELECT, SHOW, etc), the execute step can be told to fetch only a limited number of rows, so that multiple execute steps might be needed to complete the operation.

The backend can keep track of multiple prepared statements and portals (but note that these exist only within a session, and are never shared across sessions). Existing prepared statements and portals are referenced by names assigned when they were created. In addition, an “unnamed” prepared statement and portal exist. Although these behave largely the same as named objects, operations on them are optimized for the case of executing a query only once and then discarding it, whereas operations on named objects are optimized on the expectation of multiple uses.

51.1.3. Formats and Format Codes

Data of a particular data type might be transmitted in any of several different formats. As of PostgreSQL 7.4 the only supported formats are “text” and “binary”, but the protocol makes provision for future extensions. The desired format for any value is specified by a format code. Clients can specify a format code for each transmitted parameter value and for each column of a query result. Text has format code zero, binary has format code one, and all other format codes are reserved for future definition.

The text representation of values is whatever strings are produced and accepted by the input/output conversion functions for the particular data type. In the transmitted representation, there is no trailing
null character; the frontend must add one to received values if it wants to process them as C strings. (The text format does not allow embedded nulls, by the way.)

Binary representations for integers use network byte order (most significant byte first). For other data types consult the documentation or source code to learn about the binary representation. Keep in mind that binary representations for complex data types might change across server versions; the text format is usually the more portable choice.

51.2. Message Flow

This section describes the message flow and the semantics of each message type. (Details of the exact representation of each message appear in Section 51.5.) There are several different sub-protocols depending on the state of the connection: start-up, query, function call, COPY, and termination. There are also special provisions for asynchronous operations (including notification responses and command cancellation), which can occur at any time after the start-up phase.

51.2.1. Start-up

To begin a session, a frontend opens a connection to the server and sends a startup message. This message includes the names of the user and of the database the user wants to connect to; it also identifies the particular protocol version to be used. (Optionally, the startup message can include additional settings for run-time parameters.) The server then uses this information and the contents of its configuration files (such as pg_hba.conf) to determine whether the connection is provisionally acceptable, and what additional authentication is required (if any).

The server then sends an appropriate authentication request message, to which the frontend must reply with an appropriate authentication response message (such as a password). For all authentication methods except GSSAPI and SSPI, there is at most one request and one response. In some methods, no response at all is needed from the frontend, and so no authentication request occurs. For GSSAPI and SSPI, multiple exchanges of packets may be needed to complete the authentication.

The authentication cycle ends with the server either rejecting the connection attempt (ErrorResponse), or sending AuthenticationOk.

The possible messages from the server in this phase are:

ErrorResponse
The connection attempt has been rejected. The server then immediately closes the connection.

AuthenticationOk
The authentication exchange is successfully completed.

AuthenticationKerberosV5
The frontend must now take part in a Kerberos V5 authentication dialog (not described here, part of the Kerberos specification) with the server. If this is successful, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse. This is no longer supported.

AuthenticationCleartextPassword
The frontend must now send a PasswordMessage containing the password in clear-text form. If this is the correct password, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse.
AuthenticationMD5Password

The frontend must now send a PasswordMessage containing the password (with user name) encrypted via MD5, then encrypted again using the 4-byte random salt specified in the AuthenticationMD5Password message. If this is the correct password, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse. The actual PasswordMessage can be computed in SQL as `concat('md5', md5(concat(md5(concat(password, username)), random-salt)))`. (Keep in mind the `md5()` function returns its result as a hex string.)

AuthenticationSCMCredential

This response is only possible for local Unix-domain connections on platforms that support SCM credential messages. The frontend must issue an SCM credential message and then send a single data byte. (The contents of the data byte are uninteresting; it’s only used to ensure that the server waits long enough to receive the credential message.) If the credential is acceptable, the server responds with an AuthenticationOk, otherwise it responds with an ErrorResponse. (This message type is only issued by pre-9.1 servers. It may eventually be removed from the protocol specification.)

AuthenticationGSS

The frontend must now initiate a GSSAPI negotiation. The frontend will send a PasswordMessage with the first part of the GSSAPI data stream in response to this. If further messages are needed, the server will respond with AuthenticationGSSContinue.

AuthenticationSSPI

The frontend must now initiate a SSPI negotiation. The frontend will send a PasswordMessage with the first part of the SSPI data stream in response to this. If further messages are needed, the server will respond with AuthenticationGSSContinue.

AuthenticationGSSContinue

This message contains the response data from the previous step of GSSAPI or SSPI negotiation (AuthenticationGSS, AuthenticationSSPI or a previous AuthenticationGSSContinue). If the GSSAPI or SSPI data in this message indicates more data is needed to complete the authentication, the frontend must send that data as another PasswordMessage. If GSSAPI or SSPI authentication is completed by this message, the server will next send AuthenticationOk to indicate successful authentication or ErrorResponse to indicate failure.

NegotiateProtocolVersion

The server does not support the minor protocol version requested by the client, but does support an earlier version of the protocol; this message indicates the highest supported minor version. This message will also be sent if the client requested unsupported protocol options (i.e. beginning with `_pq_`) in the startup packet. This message will be followed by an ErrorResponse or a message indicating the success or failure of authentication.

If the frontend does not support the authentication method requested by the server, then it should immediately close the connection.

After having received AuthenticationOk, the frontend must wait for further messages from the server. In this phase a backend process is being started, and the frontend is just an interested bystander. It is still possible for the startup attempt to fail (ErrorResponse) or the server to decline support for the requested minor protocol version (NegotiateProtocolVersion), but in the normal case the backend will send some ParameterStatus messages, BackendKeyData, and finally ReadyForQuery.
Chapter 51. Frontend/Backend Protocol

During this phase the backend will attempt to apply any additional run-time parameter settings that were given in the startup message. If successful, these values become session defaults. An error causes ErrorResponse and exit.

The possible messages from the backend in this phase are:

**BackendKeyData**

This message provides secret-key data that the frontend must save if it wants to be able to issue cancel requests later. The frontend should not respond to this message, but should continue listening for a ReadyForQuery message.

**ParameterStatus**

This message informs the frontend about the current (initial) setting of backend parameters, such as client_encoding or DateStyle. The frontend can ignore this message, or record the settings for its future use; see Section 51.2.6 for more details. The frontend should not respond to this message, but should continue listening for a ReadyForQuery message.

**ReadyForQuery**

Start-up is completed. The frontend can now issue commands.

**ErrorResponse**

Start-up failed. The connection is closed after sending this message.

**NoticeResponse**

A warning message has been issued. The frontend should display the message but continue listening for ReadyForQuery or ErrorResponse.

The ReadyForQuery message is the same one that the backend will issue after each command cycle. Depending on the coding needs of the frontend, it is reasonable to consider ReadyForQuery as starting a command cycle, or to consider ReadyForQuery as ending the start-up phase and each subsequent command cycle.

### 51.2.2. Simple Query

A simple query cycle is initiated by the frontend sending a Query message to the backend. The message includes an SQL command (or commands) expressed as a text string. The backend then sends one or more response messages depending on the contents of the query command string, and finally a ReadyForQuery response message. ReadyForQuery informs the frontend that it can safely send a new command. (It is not actually necessary for the frontend to wait for ReadyForQuery before issuing another command, but the frontend must then take responsibility for figuring out what happens if the earlier command fails and already-issued later commands succeed.)

The possible response messages from the backend are:

**CommandComplete**

An SQL command completed normally.

**CopyInResponse**

The backend is ready to copy data from the frontend to a table; see Section 51.2.5.

**CopyOutResponse**

The backend is ready to copy data from a table to the frontend; see Section 51.2.5.
Chapter 51. Frontend/Backend Protocol

RowDescription
Indicates that rows are about to be returned in response to a `SELECT`, `FETCH`, etc query. The contents of this message describe the column layout of the rows. This will be followed by a `DataRow` message for each row being returned to the frontend.

DataRow
One of the set of rows returned by a `SELECT`, `FETCH`, etc query.

EmptyQueryResponse
An empty query string was recognized.

ErrorResponse
An error has occurred.

ReadyForQuery
Processing of the query string is complete. A separate message is sent to indicate this because the query string might contain multiple SQL commands. (CommandComplete marks the end of processing one SQL command, not the whole string.) ReadyForQuery will always be sent, whether processing terminates successfully or with an error.

NoticeResponse
A warning message has been issued in relation to the query. Notices are in addition to other responses, i.e., the backend will continue processing the command.

The response to a `SELECT` query (or other queries that return row sets, such as `EXPLAIN` or `SHOW`) normally consists of RowDescription, zero or more `DataRow` messages, and then CommandComplete. `COPY` to or from the frontend invokes special protocol as described in Section 51.2.5. All other query types normally produce only a CommandComplete message.

Since a query string could contain several queries (separated by semicolons), there might be several such response sequences before the backend finishes processing the query string. ReadyForQuery is issued when the entire string has been processed and the backend is ready to accept a new query string.

If a completely empty (no contents other than whitespace) query string is received, the response is EmptyQueryResponse followed by ReadyForQuery.

In the event of an error, ErrorResponse is issued followed by ReadyForQuery. All further processing of the query string is aborted by ErrorResponse (even if more queries remained in it). Note that this might occur partway through the sequence of messages generated by an individual query.

In simple Query mode, the format of retrieved values is always text, except when the given command is a `FETCH` from a cursor declared with the `BINARY` option. In that case, the retrieved values are in binary format. The format codes given in the RowDescription message tell which format is being used.

A frontend must be prepared to accept ErrorResponse and NoticeResponse messages whenever it is expecting any other type of message. See also Section 51.2.6 concerning messages that the backend might generate due to outside events.

Recommended practice is to code frontends in a state-machine style that will accept any message type at any time that it could make sense, rather than wiring in assumptions about the exact sequence of messages.
51.2.3. Extended Query

The extended query protocol breaks down the above-described simple query protocol into multiple steps. The results of preparatory steps can be re-used multiple times for improved efficiency. Furthermore, additional features are available, such as the possibility of supplying data values as separate parameters instead of having to insert them directly into a query string.

In the extended protocol, the frontend first sends a Parse message, which contains a textual query string, optionally some information about data types of parameter placeholders, and the name of a destination prepared-statement object (an empty string selects the unnamed prepared statement). The response is either ParseComplete or ErrorResponse. Parameter data types can be specified by OID; if not given, the parser attempts to infer the data types in the same way as it would do for untyped literal string constants.

**Note:** A parameter data type can be left unspecified by setting it to zero, or by making the array of parameter type OIDs shorter than the number of parameter symbols ($n$) used in the query string. Another special case is that a parameter's type can be specified as void (that is, the OID of the void pseudotype). This is meant to allow parameter symbols to be used for function parameters that are actually OUT parameters. Ordinarily there is no context in which a void parameter could be used, but if such a parameter symbol appears in a function's parameter list, it is effectively ignored. For example, a function call such as `foo($1,$2,$3,$4)` could match a function with two IN and two OUT arguments, if $3$ and $4$ are specified as having type void.

**Note:** The query string contained in a Parse message cannot include more than one SQL statement; else a syntax error is reported. This restriction does not exist in the simple-query protocol, but it does exist in the extended protocol, because allowing prepared statements or portals to contain multiple commands would complicate the protocol unduly.

If successfully created, a named prepared-statement object lasts till the end of the current session, unless explicitly destroyed. An unnamed prepared statement lasts only until the next Parse statement specifying the unnamed statement as destination is issued. (Note that a simple Query message also destroys the unnamed statement.) Named prepared statements must be explicitly closed before they can be redefined by another Parse message, but this is not required for the unnamed statement. Named prepared statements can also be created and accessed at the SQL command level, using PREPARE and EXECUTE.

Once a prepared statement exists, it can be readied for execution using a Bind message. The Bind message gives the name of the source prepared statement (empty string denotes the unnamed prepared statement), the name of the destination portal (empty string denotes the unnamed portal), and the values to use for any parameter placeholders present in the prepared statement. The supplied parameter set must match those needed by the prepared statement. (If you declared any void parameters in the Parse message, pass NULL values for them in the Bind message.) Bind also specifies the format to use for any data returned by the query; the format can be specified overall, or per-column. The response is either BindComplete or ErrorResponse.

**Note:** The choice between text and binary output is determined by the format codes given in Bind, regardless of the SQL command involved. The BINAR Y attribute in cursor declarations is irrelevant when using extended query protocol.
Query planning typically occurs when the Bind message is processed. If the prepared statement has
no parameters, or is executed repeatedly, the server might save the created plan and re-use it during
subsequent Bind messages for the same prepared statement. However, it will do so only if it finds that
a generic plan can be created that is not much less efficient than a plan that depends on the specific
parameter values supplied. This happens transparently so far as the protocol is concerned.

If successfully created, a named portal object lasts till the end of the current transaction, unless ex-
licitly destroyed. An unnamed portal is destroyed at the end of the transaction, or as soon as the
next Bind statement specifying the unnamed portal as destination is issued. (Note that a simple Query
message also destroys the unnamed portal.) Named portals must be explicitly closed before they can
be redefined by another Bind message, but this is not required for the unnamed portal. Named portals
can also be created and accessed at the SQL command level, using DECLARE CURSOR and FETCH.

Once a portal exists, it can be executed using an Execute message. The Execute message specifies
the portal name (empty string denotes the unnamed portal) and a maximum result-row count (zero
meaning “fetch all rows”). The result-row count is only meaningful for portals containing commands
that return row sets; in other cases the command is always executed to completion, and the row count
is ignored. The possible responses to Execute are the same as those described above for queries issued
via simple query protocol, except that Execute doesn’t cause ReadyForQuery or RowDescription to
be issued.

If Execute terminates before completing the execution of a portal (due to reaching a nonzero result-
row count), it will send a PortalSuspended message; the appearance of this message tells the frontend
that another Execute should be issued against the same portal to complete the operation. The Com-
mandComplete message indicating completion of the source SQL command is not sent until the por-
tal’s execution is completed. Therefore, an Execute phase is always terminated by the appearance of
exactly one of these messages: CommandComplete, EmptyQueryResponse (if the portal was created
from an empty query string), ErrorResponse, or PortalSuspended.

At completion of each series of extended-query messages, the frontend should issue a Sync message.
This parameterless message causes the backend to close the current transaction if it’s not inside a
BEGIN/COMMIT transaction block (“close” meaning to commit if no error, or roll back if error). Then
a ReadyForQuery response is issued. The purpose of Sync is to provide a resynchronization point for
error recovery. When an error is detected while processing any extended-query message, the backend
issues ErrorResponse, then reads and discards messages until a Sync is reached, then issues Ready-
ForQuery and returns to normal message processing. (But note that no skipping occurs if an error is
detected while processing Sync — this ensures that there is one and only one ReadyForQuery sent for
each Sync.)

*Note:* Sync does not cause a transaction block opened with BEGIN to be closed. It is possible to
detect this situation since the ReadyForQuery message includes transaction status information.

In addition to these fundamental, required operations, there are several optional operations that can
be used with extended-query protocol.

The Describe message (portal variant) specifies the name of an existing portal (or an empty string
for the unnamed portal). The response is a RowDescription message describing the rows that will be
returned by executing the portal; or a NoData message if the portal does not contain a query that will
return rows; or ErrorResponse if there is no such portal.

The Describe message (statement variant) specifies the name of an existing prepared statement (or an
empty string for the unnamed prepared statement). The response is a ParameterDescription message
describing the parameters needed by the statement, followed by a RowDescription message describing
the rows that will be returned when the statement is eventually executed (or a NoData message if the
statement will not return rows). ErrorResponse is issued if there is no such prepared statement. Note that since Bind has not yet been issued, the formats to be used for returned columns are not yet known to the backend; the format code fields in the RowDescription message will be zeroes in this case.

**Tip:** In most scenarios the frontend should issue one or the other variant of Describe before issuing Execute, to ensure that it knows how to interpret the results it will get back.

The Close message closes an existing prepared statement or portal and releases resources. It is not an error to issue Close against a nonexistent statement or portal name. The response is normally CloseComplete, but could be ErrorResponse if some difficulty is encountered while releasing resources. Note that closing a prepared statement implicitly closes any open portals that were constructed from that statement.

The Flush message does not cause any specific output to be generated, but forces the backend to deliver any data pending in its output buffers. A Flush must be sent after any extended-query command except Sync, if the frontend wishes to examine the results of that command before issuing more commands. Without Flush, messages returned by the backend will be combined into the minimum possible number of packets to minimize network overhead.

**Note:** The simple Query message is approximately equivalent to the series Parse, Bind, portal Describe, Execute, Close, Sync, using the unnamed prepared statement and portal objects and no parameters. One difference is that it will accept multiple SQL statements in the query string, automatically performing the bind/describe/execute sequence for each one in succession. Another difference is that it will not return ParseComplete, BindComplete, CloseComplete, or NoData messages.

### 51.2.4. Function Call

The Function Call sub-protocol allows the client to request a direct call of any function that exists in the database’s `pg_proc` system catalog. The client must have execute permission for the function.

**Note:** The Function Call sub-protocol is a legacy feature that is probably best avoided in new code. Similar results can be accomplished by setting up a prepared statement that does `SELECT function($1, ...)`. The Function Call cycle can then be replaced with Bind/Execute.

A Function Call cycle is initiated by the frontend sending a FunctionCall message to the backend. The backend then sends one or more response messages depending on the results of the function call, and finally a ReadyForQuery response message. ReadyForQuery informs the frontend that it can safely send a new query or function call.

The possible response messages from the backend are:

- **ErrorResponse**
  - An error has occurred.
- **FunctionCallResponse**
  - The function call was completed and returned the result given in the message. (Note that the Function Call protocol can only handle a single scalar result, not a row type or set of results.)
ReadyForQuery

Processing of the function call is complete. ReadyForQuery will always be sent, whether processing terminates successfully or with an error.

NoticeResponse

A warning message has been issued in relation to the function call. Notices are in addition to other responses, i.e., the backend will continue processing the command.

### 51.2.5. COPY Operations

The **COPY** command allows high-speed bulk data transfer to or from the server. Copy-in and copy-out operations each switch the connection into a distinct sub-protocol, which lasts until the operation is completed.

Copy-in mode (data transfer to the server) is initiated when the backend executes a **COPY FROM STDIN** SQL statement. The backend sends a CopyInResponse message to the frontend. The frontend should then send zero or more CopyData messages, forming a stream of input data. (The message boundaries are not required to have anything to do with row boundaries, although that is often a reasonable choice.) The frontend can terminate the copy-in mode by sending either a CopyDone message (allowing successful termination) or a CopyFail message (which will cause the **COPY** SQL statement to fail with an error). The backend then reverts to the command-processing mode it was in before the **COPY** started, which will be either simple or extended query protocol. It will next send either CommandComplete (if successful) or ErrorResponse (if not).

In the event of a backend-detected error during copy-in mode (including receipt of a CopyFail message), the backend will issue an ErrorResponse message. If the **COPY** command was issued via an extended-query message, the backend will now discard frontend messages until a Sync message is received, then it will issue ReadyForQuery and return to normal processing. If the **COPY** command was issued in a simple Query message, the rest of that message is discarded and ReadyForQuery is issued. In either case, any subsequent CopyData, CopyDone, or CopyFail messages issued by the frontend will simply be dropped.

The backend will ignore Flush and Sync messages received during copy-in mode. Receipt of any other non-copy message type constitutes an error that will abort the copy-in state as described above. (The exception for Flush and Sync is for the convenience of client libraries that always send Flush or Sync after an Execute message, without checking whether the command to be executed is a **COPY FROM STDIN**.)

Copy-out mode (data transfer from the server) is initiated when the backend executes a **COPY TO STDOUT** SQL statement. The backend sends a CopyOutResponse message to the frontend, followed by zero or more CopyData messages (always one per row), followed by CopyDone. The backend then reverts to the command-processing mode it was in before the **COPY** started, and sends CommandComplete. The frontend cannot abort the transfer (except by closing the connection or issuing a Cancel request), but it can discard unwanted CopyData and CopyDone messages.

In the event of a backend-detected error during copy-out mode, the backend will issue an ErrorResponse message and revert to normal processing. The frontend should treat receipt of ErrorResponse as terminating the copy-out mode.

It is possible for NoticeResponse and ParameterStatus messages to be interspersed between CopyData messages; frontends must handle these cases, and should be prepared for other asynchronous message types as well (see Section 51.2.6). Otherwise, any message type other than CopyData or CopyDone may be treated as terminating copy-out mode.
Chapter 51. Frontend/Backend Protocol

There is another Copy-related mode called copy-both, which allows high-speed bulk data transfer to and from the server. Copy-both mode is initiated when a backend in walsender mode executes a START_REPLICATION statement. The backend sends a CopyBothResponse message to the frontend. Both the backend and the frontend may then send CopyData messages until either end sends a CopyDone message. After the client sends a CopyDone message, the connection goes from copy-both mode to copy-out mode, and the client may not send any more CopyData messages. Similarly, when the server sends a CopyDone message, the connection goes into copy-in mode, and the server may not send any more CopyData messages. After both sides have sent a CopyDone message, the copy mode is terminated, and the backend reverts to the command-processing mode. In the event of a backend-detected error during copy-both mode, the backend will issue an ErrorResponse message, discard frontend messages until a Sync message is received, and then issue ReadyForQuery and return to normal processing. The frontend should treat receipt of ErrorResponse as terminating the copy in both directions; no CopyDone should be sent in this case. See Section 51.3 for more information on the subprotocol transmitted over copy-both mode.

The CopyInResponse, CopyOutResponse and CopyBothResponse messages include fields that inform the frontend of the number of columns per row and the format codes being used for each column. (As of the present implementation, all columns in a given COPY operation will use the same format, but the message design does not assume this.)

51.2.6. Asynchronous Operations

There are several cases in which the backend will send messages that are not specifically prompted by the frontend’s command stream. Frontends must be prepared to deal with these messages at any time, even when not engaged in a query. At minimum, one should check for these cases before beginning to read a query response.

It is possible for NoticeResponse messages to be generated due to outside activity; for example, if the database administrator commands a “fast” database shutdown, the backend will send a NoticeResponse indicating this fact before closing the connection. Accordingly, frontends should always be prepared to accept and display NoticeResponse messages, even when the connection is nominally idle.

ParameterStatus messages will be generated whenever the active value changes for any of the parameters the backend believes the frontend should know about. Most commonly this occurs in response to a SET SQL command executed by the frontend, and this case is effectively synchronous — but it is also possible for parameter status changes to occur because the administrator changed a configuration file and then sent the SIGHUP signal to the server. Also, if a SET command is rolled back, an appropriate ParameterStatus message will be generated to report the current effective value.

At present there is a hard-wired set of parameters for which ParameterStatus will be generated: they are server_version, server_encoding, client_encoding, application_name, is_superuser, session_authorization, DateStyle, IntervalStyle, TimeZone, integer_datetimes, and standard_conforming_strings. (server_encoding, TimeZone, and integer_datetimes were not reported by releases before 8.0; standard_conforming_strings was not reported by releases before 8.1; IntervalStyle was not reported by releases before 8.4; application_name was not reported by releases before 9.0.) Note that server_version, server_encoding and integer_datetimes are pseudo-parameters that cannot change after startup. This set might change in the future, or even become configurable. Accordingly, a frontend should simply ignore ParameterStatus for parameters that it does not understand or care about.

If a frontend issues a LISTEN command, then the backend will send a NotificationResponse message (not to be confused with NoticeResponse!) whenever a NOTIFY command is executed for the same
Note: At present, NotificationResponse can only be sent outside a transaction, and thus it will not occur in the middle of a command-response series, though it might occur just before ReadyForQuery. It is unwise to design frontend logic that assumes that, however. Good practice is to be able to accept NotificationResponse at any point in the protocol.

51.2.7. Canceling Requests in Progress

During the processing of a query, the frontend might request cancellation of the query. The cancel request is not sent directly on the open connection to the backend for reasons of implementation efficiency: we don’t want to have the backend constantly checking for new input from the frontend during query processing. Cancel requests should be relatively infrequent, so we make them slightly cumbersome in order to avoid a penalty in the normal case.

To issue a cancel request, the frontend opens a new connection to the server and sends a CancelRequest message, rather than the StartupMessage message that would ordinarily be sent across a new connection. The server will process this request and then close the connection. For security reasons, no direct reply is made to the cancel request message.

A CancelRequest message will be ignored unless it contains the same key data (PID and secret key) passed to the frontend during connection start-up. If the request matches the PID and secret key for a currently executing backend, the processing of the current query is aborted. (In the existing implementation, this is done by sending a special signal to the backend process that is processing the query.)

The cancellation signal might or might not have any effect — for example, if it arrives after the backend has finished processing the query, then it will have no effect. If the cancellation is effective, it results in the current command being terminated early with an error message.

The upshot of all this is that for reasons of both security and efficiency, the frontend has no direct way to tell whether a cancel request has succeeded. It must continue to wait for the backend to respond to the query. Issuing a cancel simply improves the odds that the current query will finish soon, and improves the odds that it will fail with an error message instead of succeeding.

Since the cancel request is sent across a new connection to the server and not across the regular frontend/backend communication link, it is possible for the cancel request to be issued by any process, not just the frontend whose query is to be canceled. This might provide additional flexibility when building multiple-process applications. It also introduces a security risk, in that unauthorized persons might try to cancel queries. The security risk is addressed by requiring a dynamically generated secret key to be supplied in cancel requests.

51.2.8. Termination

The normal, graceful termination procedure is that the frontend sends a Terminate message and immediately closes the connection. On receipt of this message, the backend closes the connection and terminates.

In rare cases (such as an administrator-commanded database shutdown) the backend might disconnect without any frontend request to do so. In such cases the backend will attempt to send an error or notice message giving the reason for the disconnection before it closes the connection.
Other termination scenarios arise from various failure cases, such as core dump at one end or the other, loss of the communications link, loss of message-boundary synchronization, etc. If either frontend or backend sees an unexpected closure of the connection, it should clean up and terminate. The frontend has the option of launching a new backend by recontacting the server if it doesn’t want to terminate itself. Closing the connection is also advisable if an unrecognizable message type is received, since this probably indicates loss of message-boundary sync.

For either normal or abnormal termination, any open transaction is rolled back, not committed. One should note however that if a frontend disconnects while a non-SELECT query is being processed, the backend will probably finish the query before noticing the disconnection. If the query is outside any transaction block (BEGIN ... COMMIT sequence) then its results might be committed before the disconnection is recognized.

51.2.9. SSL Session Encryption

If PostgreSQL was built with SSL support, frontend/backend communications can be encrypted using SSL. This provides communication security in environments where attackers might be able to capture the session traffic. For more information on encrypting PostgreSQL sessions with SSL, see Section 18.9.

To initiate an SSL-encrypted connection, the frontend initially sends an SSLRequest message rather than a StartupMessage. The server then responds with a single byte containing S or N, indicating that it is willing or unwilling to perform SSL, respectively. The frontend might close the connection at this point if it is dissatisfied with the response. To continue after S, perform an SSL startup handshake (not described here, part of the SSL specification) with the server. If this is successful, continue with sending the usual StartupMessage. In this case the StartupMessage and all subsequent data will be SSL-encrypted. To continue after N, send the usual StartupMessage and proceed without encryption.

The frontend should also be prepared to handle an ErrorMessage response to SSLRequest from the server. This would only occur if the server predates the addition of SSL support to PostgreSQL. (Such servers are now very ancient, and likely do not exist in the wild anymore.) In this case the connection must be closed, but the frontend might choose to open a fresh connection and proceed without requesting SSL.

An initial SSLRequest can also be used in a connection that is being opened to send a CancelRequest message.

While the protocol itself does not provide a way for the server to force SSL encryption, the administrator can configure the server to reject unencrypted sessions as a byproduct of authentication checking.

51.3. Streaming Replication Protocol

To initiate streaming replication, the frontend sends the replication parameter in the startup message. A Boolean value of true tells the backend to go into walsender mode, wherein a small set of replication commands can be issued instead of SQL statements. Only the simple query protocol can be used in walsender mode. Replication commands are logged in the server log when log_replication_commands is enabled. Passing database as the value instructs walsender to connect to the database specified in the dbname parameter, which will allow the connection to be used for logical replication from that database.

For the purpose of testing replication commands, you can make a replication connection via psql or any other libpq-using tool with a connection string including the replication option, e.g.:
However, it is often more useful to use `pg_receivexlog` (for physical replication) or `pg_recvlogical` (for logical replication).

The commands accepted in walsender mode are:

**IDENTIFY_SYSTEM**

Requests the server to identify itself. Server replies with a result set of a single row, containing four fields:

- **systemid** (text)
  The unique system identifier identifying the cluster. This can be used to check that the base backup used to initialize the standby came from the same cluster.

- **timeline** (int4)
  Current timeline ID. Also useful to check that the standby is consistent with the master.

- **xlogpos** (text)
  Current xlog flush location. Useful to get a known location in the transaction log where streaming can start.

- **dbname** (text)
  Database connected to or null.

**TIMELINE_HISTORY tli**

Requests the server to send over the timeline history file for timeline `tli`. Server replies with a result set of a single row, containing two fields:

- **filename** (text)
  File name of the timeline history file, e.g., `00000002.history`.

- **content** (bytea)
  Contents of the timeline history file.

**CREATE_REPLICATION_SLOT slot_name [ PHYSICAL [ RESERVE_WAL ] ] | LOGICAL output_plugin ]**

Create a physical or logical replication slot. See Section 26.2.6 for more about replication slots.

- **slot_name**
  The name of the slot to create. Must be a valid replication slot name (see Section 26.2.6.1).

- **output_plugin**
  The name of the output plugin used for logical decoding (see Section 47.6).

**RESERVE_WAL**

Specify that this physical replication slot reserves WAL immediately. Otherwise, WAL is only reserved upon connection from a streaming replication client.
Chapter 51. Frontend/Backend Protocol

START_REPLICATION [ SLOT slot_name ] [ PHYSICAL ] XXX/XXX [ TIMELINE tli ]

Instructs server to start streaming WAL, starting at WAL position XXX/XXX. If TIMELINE option is specified, streaming starts on timeline tli; otherwise, the server’s current timeline is selected. The server can reply with an error, for example if the requested section of WAL has already been recycled. On success, server responds with a CopyBothResponse message, and then starts to stream WAL to the frontend.

If a slot’s name is provided via slot_name, it will be updated as replication progresses so that the server knows which WAL segments, and if hot_standby_feedback is on which transactions, are still needed by the standby.

If the client requests a timeline that’s not the latest but is part of the history of the server, the server will stream all the WAL on that timeline starting from the requested start point up to the point where the server switched to another timeline. If the client requests streaming at exactly the end of an old timeline, the server responds immediately with CommandComplete without entering COPY mode.

After streaming all the WAL on a timeline that is not the latest one, the server will end streaming by exiting the COPY mode. When the client acknowledges this by also exiting COPY mode, the server sends a result set with one row and two columns, indicating the next timeline in this server’s history. The first column is the next timeline’s ID (type int8), and the second column is the WAL position where the switch happened (type text). Usually, the switch position is the end of the WAL that was streamed, but there are corner cases where the server can send some WAL from the old timeline that it has not itself replayed before promoting. Finally, the server sends CommandComplete message, and is ready to accept a new command.

WAL data is sent as a series of CopyData messages. (This allows other information to be intermixed; in particular the server can send an ErrorResponse message if it encounters a failure after beginning to stream.) The payload of each CopyData message from server to the client contains a message of one of the following formats:

XLogData (B)

Byte1('w')

Identifies the message as WAL data.

Int64

The starting point of the WAL data in this message.

Int64

The current end of WAL on the server.

Int64

The server’s system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Byte:

A section of the WAL data stream.

A single WAL record is never split across two XLogData messages. When a WAL record crosses a WAL page boundary, and is therefore already split using continuation records, it can be split at the page boundary. In other words, the first main WAL record and its continuation records can be sent in different XLogData messages.
Primary keepalive message (B)

Byte1('k')

Identifies the message as a sender keepalive.

Int64

The current end of WAL on the server.

Int64

The server’s system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Byte1

1 means that the client should reply to this message as soon as possible, to avoid a timeout disconnect. 0 otherwise.

The receiving process can send replies back to the sender at any time, using one of the following message formats (also in the payload of a CopyData message):

Standby status update (F)

Byte1('r')

Identifies the message as a receiver status update.

Int64

The location of the last WAL byte + 1 received and written to disk in the standby.

Int64

The location of the last WAL byte + 1 flushed to disk in the standby.

Int64

The location of the last WAL byte + 1 applied in the standby.

Int64

The client’s system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Byte1

If 1, the client requests the server to reply to this message immediately. This can be used to ping the server, to test if the connection is still healthy.

Hot Standby feedback message (F)

Byte1('h')

Identifies the message as a Hot Standby feedback message.
Int64
The client’s system clock at the time of transmission, as microseconds since midnight on 2000-01-01.

Int32
The standby’s current xmin. This may be 0, if the standby is sending notification that Hot Standby feedback will no longer be sent on this connection. Later non-zero messages may reinitiate the feedback mechanism.

Int32
The standby’s current epoch.

START_REPLICATION SLOT slot_name LOGICAL XXX/XXX [ (option_name [option_value] [, ...]) ]
Instructs server to start streaming WAL for logical replication, starting at WAL position XXX/XXX. The server can reply with an error, for example if the requested section of WAL has already been recycled. On success, server responds with a CopyBothResponse message, and then starts to stream WAL to the frontend.

The messages inside the CopyBothResponse messages are of the same format documented for \texttt{START_REPLICATION ... PHYSICAL}.

The output plugin associated with the selected slot is used to process the output for streaming.

\texttt{SLOT slot_name}
The name of the slot to stream changes from. This parameter is required, and must correspond to an existing logical replication slot created with \texttt{CREATE_REPLICATION_SLOT} in \texttt{LOGICAL} mode.

XXX/XXX
The WAL position to begin streaming at.

\texttt{option_name}
The name of an option passed to the slot’s logical decoding plugin.

\texttt{option_value}
Optional value, in the form of a string constant, associated with the specified option.

DROP_REPLICATION_SLOT slot_name
Drops a replication slot, freeing any reserved server-side resources. If the slot is currently in use by an active connection, this command fails.

\texttt{slot_name}
The name of the slot to drop.

BASE_BACKUP [LABEL ‘label’] [PROGRESS] [FAST] [WAL] [NOWAIT] [MAX_RATE rate] [TABLESPACE_MAP]
Instructs the server to start streaming a base backup. The system will automatically be put in backup mode before the backup is started, and taken out of it when the backup is complete. The
following options are accepted:

LABEL 'label'

Sets the label of the backup. If none is specified, a backup label of base backup will be used. The quoting rules for the label are the same as a standard SQL string with standard_conforming_strings turned on.

PROGRESS

Request information required to generate a progress report. This will send back an approximate size in the header of each tablespace, which can be used to calculate how far along the stream is done. This is calculated by enumerating all the file sizes once before the transfer is even started, and might as such have a negative impact on the performance. In particular, it might take longer before the first data is streamed. Since the database files can change during the backup, the size is only approximate and might both grow and shrink between the time of approximation and the sending of the actual files.

FAST

Request a fast checkpoint.

WAL

Include the necessary WAL segments in the backup. This will include all the files between start and stop backup in the pg_xlog directory of the base directory tar file.

NOWAIT

By default, the backup will wait until the last required WAL segment has been archived, or emit a warning if log archiving is not enabled. Specifying NOWAIT disables both the waiting and the warning, leaving the client responsible for ensuring the required log is available.

MAX_RATE rate

Limit (throttle) the maximum amount of data transferred from server to client per unit of time. The expected unit is kilobytes per second. If this option is specified, the value must either be equal to zero or it must fall within the range from 32 kB through 1 GB (inclusive). If zero is passed or the option is not specified, no restriction is imposed on the transfer.

TABLESPACE_MAP

Include information about symbolic links present in the directory pg_tblspc in a file named tablespace_map. The tablespace map file includes each symbolic link name as it exists in the directory pg_tblspc/ and the full path of that symbolic link.

When the backup is started, the server will first send two ordinary result sets, followed by one or more CopyResponse results.

The first ordinary result set contains the starting position of the backup, in a single row with two columns. The first column contains the start position given in XLogRecPtr format, and the second column contains the corresponding timeline ID.

The second ordinary result set has one row for each tablespace. The fields in this row are:

spcoid(oid)

The OID of the tablespace, or null if it’s the base directory.

spclocation(text)

The full path of the tablespace directory, or null if it’s the base directory.
size (int8)

The approximate size of the tablespace, if progress report has been requested; otherwise it’s null.

After the second regular result set, one or more CopyResponse results will be sent, one for the main data directory and one for each additional tablespace other than pg_default and pg_global. The data in the CopyResponse results will be a tar format (following the “ustar interchange format” specified in the POSIX 1003.1-2008 standard) dump of the tablespace contents, except that the two trailing blocks of zeroes specified in the standard are omitted. After the tar data is complete, a final ordinary result set will be sent, containing the WAL end position of the backup, in the same format as the start position.

The tar archive for the data directory and each tablespace will contain all files in the directories, regardless of whether they are PostgreSQL files or other files added to the same directory. The only excluded files are:

- postmaster.pid
- postmaster.opts
- various temporary files created during the operation of the PostgreSQL server
- pg_xlog, including subdirectories. If the backup is run with WAL files included, a synthesized version of pg_xlog will be included, but it will only contain the files necessary for the backup to work, not the rest of the contents.
- pg_replslot is copied as an empty directory.
- Files other than regular files and directories, such as symbolic links and special device files, are skipped. (Symbolic links in pg_tblspc are maintained.)

Owner, group, and file mode are set if the underlying file system on the server supports it.

### 51.4. Message Data Types

This section describes the base data types used in messages.

**Intn(i)**

An n-bit integer in network byte order (most significant byte first). If i is specified it is the exact value that will appear, otherwise the value is variable. Eg. Int16, Int32(42).

**Intn[k]**

An array of k n-bit integers, each in network byte order. The array length k is always determined by an earlier field in the message. Eg. Int16[M].

**String(s)**

A null-terminated string (C-style string). There is no specific length limitation on strings. If s is specified it is the exact value that will appear, otherwise the value is variable. Eg. String, String("user").

**Note:** There is no predefined limit on the length of a string that can be returned by the backend. Good coding strategy for a frontend is to use an expandable buffer so that anything that fits in memory can be accepted. If that’s not feasible, read the full string and discard trailing characters that don’t fit into your fixed-size buffer.
Byte\(n(c)\)

Exactly \(n\) bytes. If the field width \(n\) is not a constant, it is always determinable from an earlier field in the message. If \(c\) is specified it is the exact value. Eg. Byte2, Byte1(\'n\').

51.5. Message Formats

This section describes the detailed format of each message. Each is marked to indicate that it can be sent by a frontend (F), a backend (B), or both (F & B). Notice that although each message includes a byte count at the beginning, the message format is defined so that the message end can be found without reference to the byte count. This aids validity checking. (The CopyData message is an exception, because it forms part of a data stream; the contents of any individual CopyData message cannot be interpretable on their own.)

AuthenticationOk (B)

Byte1(‘R’)

Identifies the message as an authentication request.

Int32(8)

Length of message contents in bytes, including self.

Int32(0)

Specifies that the authentication was successful.

AuthenticationKerberosV5 (B)

Byte1(‘R’)

Identifies the message as an authentication request.

Int32(8)

Length of message contents in bytes, including self.

Int32(2)

Specifies that Kerberos V5 authentication is required.

AuthenticationCleartextPassword (B)

Byte1(‘R’)

Identifies the message as an authentication request.

Int32(8)

Length of message contents in bytes, including self.
Chapter 51. Frontend/Backend Protocol

Int32(3)
Specifies that a clear-text password is required.

AuthenticationMD5Password (B)

Byte1('R')
Identifies the message as an authentication request.
Int32(12)
Length of message contents in bytes, including self.
Int32(5)
Specifies that an MD5-encrypted password is required.
Byte4
The salt to use when encrypting the password.

AuthenticationSCMCredential (B)

Byte1('R')
Identifies the message as an authentication request.
Int32(8)
Length of message contents in bytes, including self.
Int32(6)
Specifies that an SCM credentials message is required.

AuthenticationGSS (B)

Byte1('R')
Identifies the message as an authentication request.
Int32(8)
Length of message contents in bytes, including self.
Int32(7)
Specifies that GSSAPI authentication is required.

AuthenticationSSPI (B)

Byte1('R')
Identifies the message as an authentication request.
Int32(8)
Length of message contents in bytes, including self.
Chapter 51. Frontend/Backend Protocol

Int32(9)
   Specifies that SSPI authentication is required.

AuthenticationGSSContinue (B)

Byte1('R')
   Identifies the message as an authentication request.

Int32
   Length of message contents in bytes, including self.

Int32(8)
   Specifies that this message contains GSSAPI or SSPI data.

Byte
   GSSAPI or SSPI authentication data.

BackendKeyData (B)

Byte1('K')
   Identifies the message as cancellation key data. The frontend must save these values if it wishes to be able to issue CancelRequest messages later.

Int32(12)
   Length of message contents in bytes, including self.

Int32
   The process ID of this backend.

Int32
   The secret key of this backend.

Bind (F)

Byte1('B')
   Identifies the message as a Bind command.

Int32
   Length of message contents in bytes, including self.

String
   The name of the destination portal (an empty string selects the unnamed portal).

String
   The name of the source prepared statement (an empty string selects the unnamed prepared statement).
Chapter 51. Frontend/Backend Protocol

Int16
The number of parameter format codes that follow (denoted $c$ below). This can be zero to indicate that there are no parameters or that the parameters all use the default format (text); or one, in which case the specified format code is applied to all parameters; or it can equal the actual number of parameters.

Int16[$c$]
The parameter format codes. Each must presently be zero (text) or one (binary).

Int16
The number of parameter values that follow (possibly zero). This must match the number of parameters needed by the query.

Next, the following pair of fields appear for each parameter:

Int32
The length of the parameter value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL parameter value. No value bytes follow in the NULL case.

Byte $n$
The value of the parameter, in the format indicated by the associated format code. $n$ is the above length.

After the last parameter, the following fields appear:

Int16
The number of result-column format codes that follow (denoted $r$ below). This can be zero to indicate that there are no result columns or that the result columns should all use the default format (text); or one, in which case the specified format code is applied to all result columns (if any); or it can equal the actual number of result columns of the query.

Int16[$r$]
The result-column format codes. Each must presently be zero (text) or one (binary).

BindComplete (B)

Byte1(‘2’)
Identifies the message as a Bind-complete indicator.

Int32(4)
Length of message contents in bytes, including self.

CancelRequest (F)

Int32(16)
Length of message contents in bytes, including self.
Int32(80877102)

The cancel request code. The value is chosen to contain 1234 in the most significant 16 bits, and 5678 in the least significant 16 bits. (To avoid confusion, this code must not be the same as any protocol version number.)

Int32

The process ID of the target backend.

Int32

The secret key for the target backend.

Close (F)

Byte1(‘C’)

Identifies the message as a Close command.

Int32

Length of message contents in bytes, including self.

Byte1

‘s’ to close a prepared statement; or ‘p’ to close a portal.

String

The name of the prepared statement or portal to close (an empty string selects the unnamed prepared statement or portal).

CloseComplete (B)

Byte1(‘3’)

Identifies the message as a Close-complete indicator.

Int32(4)

Length of message contents in bytes, including self.

CommandComplete (B)

Byte1(‘C’)

Identifies the message as a command-completed response.

Int32

Length of message contents in bytes, including self.

String

The command tag. This is usually a single word that identifies which SQL command was completed.

For an INSERT command, the tag is INSERT oid rows, where rows is the number of rows inserted. oid is the object ID of the inserted row if rows is 1 and the target table has OIDs; otherwise oid is 0.
For a **DELETE** command, the tag is **DELETE rows** where **rows** is the number of rows deleted.

For an **UPDATE** command, the tag is **UPDATE rows** where **rows** is the number of rows updated.

For a **SELECT** or **CREATE TABLE AS** command, the tag is **SELECT rows** where **rows** is the number of rows retrieved.

For a **MOVE** command, the tag is **MOVE rows** where **rows** is the number of rows the cursor’s position has been changed by.

For a **FETCH** command, the tag is **FETCH rows** where **rows** is the number of rows that have been retrieved from the cursor.

For a **COPY** command, the tag is **COPY rows** where **rows** is the number of rows copied.

(Note: the row count appears only in PostgreSQL 8.2 and later.)

**CopyData (F & B)**

Byte1(‘d’)

Identifies the message as **COPY** data.

Int32

Length of message contents in bytes, including self.

Byte

Data that forms part of a **COPY** data stream. Messages sent from the backend will always correspond to single data rows, but messages sent by frontends might divide the data stream arbitrarily.

**CopyDone (F & B)**

Byte1(‘c’)

Identifies the message as a **COPY**-complete indicator.

Int32(4)

Length of message contents in bytes, including self.

**CopyFail (F)**

Byte1(‘f’)

Identifies the message as a **COPY**-failure indicator.

Int32

Length of message contents in bytes, including self.

String

An error message to report as the cause of failure.
CopyInResponse (B)

Byte1('G')
Identifies the message as a Start Copy In response. The frontend must now send copy-in data (if not prepared to do so, send a CopyFail message).

Int32
Length of message contents in bytes, including self.

Int8
0 indicates the overall COPY format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary (similar to DataRow format). See COPY for more information.

Int16
The number of columns in the data to be copied (denoted $N$ below).

Int16[$N$]
The format codes to be used for each column. Each must presently be zero (text) or one (binary). All must be zero if the overall copy format is textual.

CopyOutResponse (B)

Byte1('H')
Identifies the message as a Start Copy Out response. This message will be followed by copy-out data.

Int32
Length of message contents in bytes, including self.

Int8
0 indicates the overall COPY format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary (similar to DataRow format). See COPY for more information.

Int16
The number of columns in the data to be copied (denoted $N$ below).

Int16[$N$]
The format codes to be used for each column. Each must presently be zero (text) or one (binary). All must be zero if the overall copy format is textual.

CopyBothResponse (B)

Byte1('W')
Identifies the message as a Start Copy Both response. This message is used only for Streaming Replication.

Int32
Length of message contents in bytes, including self.
Int8
0 indicates the overall COPY format is textual (rows separated by newlines, columns separated by separator characters, etc). 1 indicates the overall copy format is binary (similar to DataRow format). See COPY for more information.

Int16
The number of columns in the data to be copied (denoted \(N\) below).

Int16[\(N\)]
The format codes to be used for each column. Each must presently be zero (text) or one (binary). All must be zero if the overall copy format is textual.

DataRow (B)

Byte1('D')
Identifies the message as a data row.

Int32
Length of message contents in bytes, including self.

Int16
The number of column values that follow (possibly zero).

Next, the following pair of fields appear for each column:

Int32
The length of the column value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL column value. No value bytes follow in the NULL case.

Byte\(n\)
The value of the column, in the format indicated by the associated format code. \(n\) is the above length.

Describe (F)

Byte1('D')
Identifies the message as a Describe command.

Int32
Length of message contents in bytes, including self.

Byte1
's' to describe a prepared statement; or 'p' to describe a portal.

String
The name of the prepared statement or portal to describe (an empty string selects the unnamed prepared statement or portal).
EmptyQueryResponse (B)

Byte1('I')
Identifies the message as a response to an empty query string. (This substitutes for CommandComplete.)

Int32(4)
Length of message contents in bytes, including self.

ErrorResponse (B)

Byte1('E')
Identifies the message as an error.

Int32
Length of message contents in bytes, including self.

The message body consists of one or more identified fields, followed by a zero byte as a terminator. Fields can appear in any order. For each field there is the following:

Byte1
A code identifying the field type; if zero, this is the message terminator and no string follows. The presently defined field types are listed in Section 51.6. Since more field types might be added in future, frontends should silently ignore fields of unrecognized type.

String
The field value.

Execute (F)

Byte1('E')
Identifies the message as an Execute command.

Int32
Length of message contents in bytes, including self.

String
The name of the portal to execute (an empty string selects the unnamed portal).

Int32
Maximum number of rows to return, if portal contains a query that returns rows (ignored otherwise). Zero denotes “no limit”.

Flush (F)

Byte1('H')
Identifies the message as a Flush command.
Int32(4)
Length of message contents in bytes, including self.

FunctionCall (F)

Byte1('F')
Identifies the message as a function call.
Int32
Length of message contents in bytes, including self.
Int32
Specifies the object ID of the function to call.
Int16
The number of argument format codes that follow (denoted \( c \) below). This can be zero to indicate that there are no arguments or that the arguments all use the default format (text); or one, in which case the specified format code is applied to all arguments; or it can equal the actual number of arguments.
Int16[\( c \)]
The argument format codes. Each must presently be zero (text) or one (binary).
Int16
Specifies the number of arguments being supplied to the function.
Next, the following pair of fields appear for each argument:
Int32
The length of the argument value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL argument value. No value bytes follow in the NULL case.
Byte\( n \)
The value of the argument, in the format indicated by the associated format code. \( n \) is the above length.
After the last argument, the following field appears:
Int16
The format code for the function result. Must presently be zero (text) or one (binary).

FunctionCallResponse (B)

Byte1('V')
Identifies the message as a function call result.
Int32
Length of message contents in bytes, including self.
Int32
The length of the function result value, in bytes (this count does not include itself). Can be zero. As a special case, -1 indicates a NULL function result. No value bytes follow in the NULL case.

Byte n
The value of the function result, in the format indicated by the associated format code. \( n \) is the above length.

NegotiateProtocolVersion (B)

Byte1('v')
Identifies the message as a protocol version negotiation message.

Int32
Length of message contents in bytes, including self.

Int32
Newest minor protocol version supported by the server for the major protocol version requested by the client.

Int32
Number of protocol options not recognized by the server.

Then, for protocol option not recognized by the server, there is the following:

String
The option name.

NoData (B)

Byte1('n')
Identifies the message as a no-data indicator.

Int32(4)
Length of message contents in bytes, including self.

NoticeResponse (B)

Byte1('N')
Identifies the message as a notice.

Int32
Length of message contents in bytes, including self.

The message body consists of one or more identified fields, followed by a zero byte as a terminator. Fields can appear in any order. For each field there is the following:
Byte1
A code identifying the field type; if zero, this is the message terminator and no string follows. The presently defined field types are listed in Section 51.6. Since more field types might be added in future, frontends should silently ignore fields of unrecognized type.

String
The field value.

NotificationResponse (B)

Byte1('A')
Identifies the message as a notification response.

Int32
Length of message contents in bytes, including self.

Int32
The process ID of the notifying backend process.

String
The name of the channel that the notify has been raised on.

String
The “payload” string passed from the notifying process.

ParameterDescription (B)

Byte1('t')
Identifies the message as a parameter description.

Int32
Length of message contents in bytes, including self.

Int16
The number of parameters used by the statement (can be zero).
Then, for each parameter, there is the following:

Int32
Specifies the object ID of the parameter data type.

ParameterStatus (B)

Byte1('S')
Identifies the message as a run-time parameter status report.

Int32
Length of message contents in bytes, including self.
String
   The name of the run-time parameter being reported.

String
   The current value of the parameter.

Parse (F)

Byte1(‘P’)
   Identifies the message as a Parse command.

Int32
   Length of message contents in bytes, including self.

String
   The name of the destination prepared statement (an empty string selects the unnamed pre-
   pared statement).

String
   The query string to be parsed.

Int16
   The number of parameter data types specified (can be zero). Note that this is not an indi-
   cation of the number of parameters that might appear in the query string, only the number
   that the frontend wants to prespecify types for.

Then, for each parameter, there is the following:

Int32
   Specifies the object ID of the parameter data type. Placing a zero here is equivalent to
   leaving the type unspecified.

ParseComplete (B)

Byte1(‘1’)
   Identifies the message as a Parse-complete indicator.

Int32(4)
   Length of message contents in bytes, including self.

PasswordMessage (F)

Byte1(‘p’)
   Identifies the message as a password response. Note that this is also used for GSSAPI and
   SSPI response messages (which is really a design error, since the contained data is not a
   null-terminated string in that case, but can be arbitrary binary data).

Int32
   Length of message contents in bytes, including self.
Chapter 51. Frontend/Backend Protocol

String
The password (encrypted, if requested).

PortalSuspended (B)

Byte1('s')
Identifies the message as a portal-suspended indicator. Note this only appears if an Execute message's row-count limit was reached.
Int32(4)
Length of message contents in bytes, including self.

Query (F)

Byte1('Q')
Identifies the message as a simple query.
Int32
Length of message contents in bytes, including self.
String
The query string itself.

ReadyForQuery (B)

Byte1('Z')
Identifies the message type. ReadyForQuery is sent whenever the backend is ready for a new query cycle.
Int32(5)
Length of message contents in bytes, including self.
Byte1
Current backend transaction status indicator. Possible values are 'I' if idle (not in a transaction block); 'T' if in a transaction block; or 'E' if in a failed transaction block (queries will be rejected until block is ended).

RowDescription (B)

Byte1('T')
Identifies the message as a row description.
Int32
Length of message contents in bytes, including self.
Int16

Specifies the number of fields in a row (can be zero).

Then, for each field, there is the following:

String

The field name.

Int32

If the field can be identified as a column of a specific table, the object ID of the table; otherwise zero.

Int16

If the field can be identified as a column of a specific table, the attribute number of the column; otherwise zero.

Int32

The object ID of the field’s data type.

Int16

The data type size (see pg_type.typlen). Note that negative values denote variable-width types.

Int32

The type modifier (see pg_attribute.atttypmod). The meaning of the modifier is type-specific.

Int16

The format code being used for the field. Currently will be zero (text) or one (binary). In a RowDescription returned from the statement variant of Describe, the format code is not yet known and will always be zero.

SSLRequest (F)

Int32(8)

Length of message contents in bytes, including self.

Int32(80877103)

The SSL request code. The value is chosen to contain 1234 in the most significant 16 bits, and 5679 in the least significant 16 bits. (To avoid confusion, this code must not be the same as any protocol version number.)

StartupMessage (F)

Int32

Length of message contents in bytes, including self.

Int32(196608)

The protocol version number. The most significant 16 bits are the major version number (3 for the protocol described here). The least significant 16 bits are the minor version number (0 for the protocol described here).
Chapter 51. Frontend/Backend Protocol

The protocol version number is followed by one or more pairs of parameter name and value strings. A zero byte is required as a terminator after the last name/value pair. Parameters can appear in any order. user is required, others are optional. Each parameter is specified as:

**String**

The parameter name. Currently recognized names are:

user

The database user name to connect as. Required; there is no default.

database

The database to connect to. Defaults to the user name.

options

Command-line arguments for the backend. (This is deprecated in favor of setting individual run-time parameters.) Spaces within this string are considered to separate arguments, unless escaped with a backslash (\); write `\` to represent a literal backslash.

In addition to the above, other parameters may be listed. Parameter names beginning with `_pq_` are reserved for use as protocol extensions, while others are treated as run-time parameters to be set at backend start time. Such settings will be applied during backend start (after parsing the command-line arguments if any) and will act as session defaults.

**String**

The parameter value.

**Sync (F)**

Byte1(‘S’) Identifies the message as a Sync command.

Int32(4) Length of message contents in bytes, including self.

**Terminate (F)**

Byte1(‘X’) Identifies the message as a termination.

Int32(4) Length of message contents in bytes, including self.

51.6. Error and Notice Message Fields

This section describes the fields that can appear in ErrorResponse and NoticeResponse messages. Each field type has a single-byte identification token. Note that any given field type should appear at most once per message.
Severity: the field contents are ERROR, FATAL, or PANIC (in an error message), or WARNING, NOTICE, DEBUG, INFO, or LOG (in a notice message), or a localized translation of one of these. Always present.

Severity: the field contents are ERROR, FATAL, or PANIC (in an error message), or WARNING, NOTICE, DEBUG, INFO, or LOG (in a notice message). This is identical to the S field except that the contents are never localized. This is present only in messages generated by PostgreSQL versions 9.6 and later.

Code: the SQLSTATE code for the error (see Appendix A). Not localizable. Always present.

Message: the primary human-readable error message. This should be accurate but terse (typically one line). Always present.

Detail: an optional secondary error message carrying more detail about the problem. Might run to multiple lines.

Hint: an optional suggestion what to do about the problem. This is intended to differ from Detail in that it offers advice (potentially inappropriate) rather than hard facts. Might run to multiple lines.

Position: the field value is a decimal ASCII integer, indicating an error cursor position as an index into the original query string. The first character has index 1, and positions are measured in characters not bytes.

Internal position: this is defined the same as the P field, but it is used when the cursor position refers to an internally generated command rather than the one submitted by the client. The q field will always appear when this field appears.

Internal query: the text of a failed internally-generated command. This could be, for example, a SQL query issued by a PL/pgSQL function.

Where: an indication of the context in which the error occurred. Presently this includes a call stack traceback of active procedural language functions and internally-generated queries. The trace is one entry per line, most recent first.

Schema name: if the error was associated with a specific database object, the name of the schema containing that object, if any.

Table name: if the error was associated with a specific table, the name of the table. (Refer to the schema name field for the name of the table’s schema.)
Chapter 51. Frontend/Backend Protocol

Column name: if the error was associated with a specific table column, the name of the column. (Refer to the schema and table name fields to identify the table.)

Data type name: if the error was associated with a specific data type, the name of the data type. (Refer to the schema name field for the name of the data type’s schema.)

Constraint name: if the error was associated with a specific constraint, the name of the constraint. Refer to fields listed above for the associated table or domain. (For this purpose, indexes are treated as constraints, even if they weren’t created with constraint syntax.)

File: the file name of the source-code location where the error was reported.

Line: the line number of the source-code location where the error was reported.

Routine: the name of the source-code routine reporting the error.

Note: The fields for schema name, table name, column name, data type name, and constraint name are supplied only for a limited number of error types; see Appendix A. Frontends should not assume that the presence of any of these fields guarantees the presence of another field. Core error sources observe the interrelationships noted above, but user-defined functions may use these fields in other ways. In the same vein, clients should not assume that these fields denote contemporary objects in the current database.

The client is responsible for formatting displayed information to meet its needs; in particular it should break long lines as needed. Newline characters appearing in the error message fields should be treated as paragraph breaks, not line breaks.

51.7. Summary of Changes since Protocol 2.0

This section provides a quick checklist of changes, for the benefit of developers trying to update existing client libraries to protocol 3.0.

The initial startup packet uses a flexible list-of-strings format instead of a fixed format. Notice that session default values for run-time parameters can now be specified directly in the startup packet. (Actually, you could do that before using the options field, but given the limited width of options and the lack of any way to quote whitespace in the values, it wasn’t a very safe technique.)

All messages now have a length count immediately following the message type byte (except for startup packets, which have no type byte). Also note that PasswordMessage now has a type byte.

ErrorResponse and NoticeResponse (‘E’ and ‘N’) messages now contain multiple fields, from which the client code can assemble an error message of the desired level of verbosity. Note that individual fields will typically not end with a newline, whereas the single string sent in the older protocol always did.

The ReadyForQuery (‘Z’) message includes a transaction status indicator.
The distinction between BinaryRow and DataRow message types is gone; the single DataRow message type serves for returning data in all formats. Note that the layout of DataRow has changed to make it easier to parse. Also, the representation of binary values has changed: it is no longer directly tied to the server’s internal representation.

There is a new “extended query” sub-protocol, which adds the frontend message types Parse, Bind, Execute, Describe, Close, Flush, and Sync, and the backend message types ParseComplete, BindComplete, PortalSuspended, ParameterDescription, NoData, and CloseComplete. Existing clients do not have to concern themselves with this sub-protocol, but making use of it might allow improvements in performance or functionality.

COPY data is now encapsulated into CopyData and CopyDone messages. There is a well-defined way to recover from errors during COPY. The special “\.” last line is not needed anymore, and is not sent during COPY OUT. (It is still recognized as a terminator during COPY IN, but its use is deprecated and will eventually be removed.) Binary COPY is supported. The CopyInResponse and CopyOutResponse messages include fields indicating the number of columns and the format of each column.

The layout of FunctionCall and FunctionCallResponse messages has changed. FunctionCall can now support passing NULL arguments to functions. It also can handle passing parameters and retrieving results in either text or binary format. There is no longer any reason to consider FunctionCall a potential security hole, since it does not offer direct access to internal server data representations.

The backend sends ParameterStatus (‘$’) messages during connection startup for all parameters it considers interesting to the client library. Subsequently, a ParameterStatus message is sent whenever the active value changes for any of these parameters.

The RowDescription (‘$’) message carries new table OID and column number fields for each column of the described row. It also shows the format code for each column.

The CursorResponse (‘$’) message is no longer generated by the backend.

The NotificationResponse (‘$’) message has an additional string field, which can carry a “payload” string passed from the NOTIFY event sender.

The EmptyQueryResponse (‘$’) message used to include an empty string parameter; this has been removed.
Chapter 52. PostgreSQL Coding Conventions

52.1. Formatting

Source code formatting uses 4 column tab spacing, with tabs preserved (i.e., tabs are not expanded to spaces). Each logical indentation level is one additional tab stop.

Layout rules (brace positioning, etc) follow BSD conventions. In particular, curly braces for the controlled blocks of `if`, `while`, `switch`, etc go on their own lines.

Limit line lengths so that the code is readable in an 80-column window. (This doesn’t mean that you must never go past 80 columns. For instance, breaking a long error message string in arbitrary places just to keep the code within 80 columns is probably not a net gain in readability.)

Do not use C++ style comments (`//` comments). Strict ANSI C compilers do not accept them. For the same reason, do not use C++ extensions such as declaring new variables mid-block.

The preferred style for multi-line comment blocks is

```c
/*
 * comment text begins here
 * and continues here
 */
```

Note that comment blocks that begin in column 1 will be preserved as-is by pgindent, but it will re-flow indented comment blocks as though they were plain text. If you want to preserve the line breaks in an indented block, add dashes like this:

```c
-------------
 * comment text begins here
 * and continues here
 *-------------
 * /
```

While submitted patches do not absolutely have to follow these formatting rules, it’s a good idea to do so. Your code will get run through pgindent before the next release, so there’s no point in making it look nice under some other set of formatting conventions. A good rule of thumb for patches is “make the new code look like the existing code around it”.

The `src/tools` directory contains sample settings files that can be used with the emacs, xemacs or vim editors to help ensure that they format code according to these conventions.

The text browsing tools more and less can be invoked as:

```bash
more -x4
less -x4
```

to make them show tabs appropriately.
52.2. Reporting Errors Within the Server

Error, warning, and log messages generated within the server code should be created using `ereport`, or its older cousin `elog`. The use of this function is complex enough to require some explanation.

There are two required elements for every message: a severity level (ranging from `DEBUG` to `PANIC`) and a primary message text. In addition there are optional elements, the most common of which is an error identifier code that follows the SQL spec’s SQLSTATE conventions. `ereport` itself is just a shell function, that exists mainly for the syntactic convenience of making message generation look like a function call in the C source code. The only parameter accepted directly by `ereport` is the severity level. The primary message text and any optional message elements are generated by calling auxiliary functions, such as `errmsg`, within the `ereport` call.

A typical call to `ereport` might look like this:

```c
ereport(ERROR,
    (errcode(ERRCODE_DIVISION_BY_ZERO),
    errmsg("division by zero")));
```

This specifies error severity level `ERROR` (a run-of-the-mill error). The `errcode` call specifies the SQLSTATE error code using a macro defined in `src/include/utils/errcodes.h`. The `errmsg` call provides the primary message text. Notice the extra set of parentheses surrounding the auxiliary function calls — these are annoying but syntactically necessary.

Here is a more complex example:

```c
ereport(ERROR,
    (errcode(ERRCODE_AMBIGUOUS_FUNCTION),
    errmsg("function %s is not unique",
        func_signature_string(funcname, nargs,
            NIL, actual_arg_types)),
    errhint("Unable to choose a best candidate function. "
        "You might need to add explicit typecasts.")));
```

This illustrates the use of format codes to embed run-time values into a message text. Also, an optional "hint" message is provided.

If the severity level is `ERROR` or higher, `ereport` aborts the execution of the user-defined function and does not return to the caller. If the severity level is lower than `ERROR`, `ereport` returns normally.

The available auxiliary routines for `ereport` are:

- `errcode(sqlerrcode)` specifies the SQLSTATE error identifier code for the condition. If this routine is not called, the error identifier defaults to `ERCODE_INTERNAL_ERROR` when the error severity level is `ERROR` or higher, `ERCODE_WARNING` when the error level is `WARNING`, otherwise (for `NOTICE` and below) `ERCODE_SUCCESSFUL_COMPLETION`. While these defaults are often convenient, always think whether they are appropriate before omitting the `errcode()` call.

- `errmsg(const char *msg, ...)` specifies the primary error message text, and possibly run-time values to insert into it. Insertions are specified by `sprintf`-style format codes. In addition to the standard format codes accepted by `sprintf`, the format code `%m` can be used to insert the error message returned by `strerror` for the current value of `errno`.¹ `%m` does not require any corresponding entry in the parameter list for `errmsg`. Note that the message string will be run through `gettext` for possible localization before format codes are processed.

---

¹. That is, the value that was current when the `ereport` call was reached; changes of `errno` within the auxiliary reporting routines will not affect it. That would not be true if you were to write `strerror` explicitly in `errmsg`'s parameter list; accordingly, do not do so.
• `errmsg_internal(const char *msg, ...)` is the same as `errmsg`, except that the message string will not be translated nor included in the internationalization message dictionary. This should be used for “cannot happen” cases that are probably not worth expending translation effort on.

• `errmsg_plural(const char *fmt_singular, const char *fmt_plural, unsigned long n, ...)`) is like `errmsg`, but with support for various plural forms of the message. `fmt_singular` is the English singular format, `fmt_plural` is the English plural format, `n` is the integer value that determines which plural form is needed, and the remaining arguments are formatted according to the selected format string. For more information see Section 53.2.2.

• `errdetail(const char *msg, ...)`) supplies an optional “detail” message; this is to be used when there is additional information that seems inappropriate to put in the primary message. The message string is processed in just the same way as for `errmsg`.

• `errdetail_internal(const char *msg, ...)`) is the same as `errdetail`, except that the message string will not be translated nor included in the internationalization message dictionary. This should be used for detail messages that are not worth expending translation effort on, for instance because they are too technical to be useful to most users.

• `errdetail_plural(const char *fmt_singular, const char *fmt_plural, unsigned long n, ...)`) is like `errdetail`, but with support for various plural forms of the message. For more information see Section 53.2.2.

• `errdetail_log(const char *msg, ...)`) is the same as `errdetail` except that this string goes only to the server log, never to the client. If both `errdetail` (or one of its equivalents above) and `errdetail_log` are used then one string goes to the client and the other to the log. This is useful for error details that are too security-sensitive or too bulky to include in the report sent to the client.

• `errdetail_log_plural(const char *fmt_singular, const char *fmt_plural, unsigned long n, ...)`) is like `errdetail_log`, but with support for various plural forms of the message. For more information see Section 53.2.2.

• `errhint(const char *msg, ...)`) supplies an optional “hint” message; this is to be used when offering suggestions about how to fix the problem, as opposed to factual details about what went wrong. The message string is processed in just the same way as for `errmsg`.

• `errcontext(const char *msg, ...)`) is not normally called directly from an `ereport` message site; rather it is used in `error_context_stack` callback functions to provide information about the context in which an error occurred, such as the current location in a PL function. The message string is processed in just the same way as for `errmsg`. Unlike the other auxiliary functions, this can be called more than once per `ereport` call; the successive strings thus supplied are concatenated with separating newlines.

• `errposition(int cursorpos)` specifies the textual location of an error within a query string. Currently it is only useful for errors detected in the lexical and syntactic analysis phases of query processing.

• `errtable(Relation rel)` specifies a relation whose name and schema name should be included as auxiliary fields in the error report.

• `errtablecol(Relation rel, int attnum)` specifies a column whose name, table name, and schema name should be included as auxiliary fields in the error report.

• `errtableconstraint(Relation rel, const char *conname)` specifies a table constraint whose name, table name, and schema name should be included as auxiliary fields in the error report. Indexes should be considered to be constraints for this purpose, whether or not they have an associated `pg_constraint` entry. Be careful to pass the underlying heap relation, not the index itself, as `rel`. 

Chapter 52. PostgreSQL Coding Conventions
Chapter 52. PostgreSQL Coding Conventions

- `errdatatype(Oid datatypeOid)` specifies a data type whose name and schema name should be included as auxiliary fields in the error report.
- `errdomainconstraint(Oid datatypeOid, const char *conname)` specifies a domain constraint whose name, domain name, and schema name should be included as auxiliary fields in the error report.
- `errcode_for_file_access()` is a convenience function that selects an appropriate SQLSTATE error identifier for a failure in a file-access-related system call. It uses the saved `errno` to determine which error code to generate. Usually this should be used in combination with `%m` in the primary error message text.
- `errcode_for_socket_access()` is a convenience function that selects an appropriate SQLSTATE error identifier for a failure in a socket-related system call.
- `errhidestmt(bool hide_stmt)` can be called to specify suppression of the STATEMENT: portion of a message in the postmaster log. Generally this is appropriate if the message text includes the current statement already.
- `errhidecontext(bool hide_ctx)` can be called to specify suppression of the CONTEXT: portion of a message in the postmaster log. This should only be used for verbose debugging messages where the repeated inclusion of context would bloat the log volume too much.

Note: At most one of the functions `errtable`, `errtablecol`, `errtableconstraint`, `errdatatype`, or `errdomainconstraint` should be used in an `ereport` call. These functions exist to allow applications to extract the name of a database object associated with the error condition without having to examine the potentially-localized error message text. These functions should be used in error reports for which it's likely that applications would wish to have automatic error handling. As of PostgreSQL 9.3, complete coverage exists only for errors in SQLSTATE class 23 (integrity constraint violation), but this is likely to be expanded in future.

There is an older function `elog` that is still heavily used. An `elog` call:

```c
elog(level, "format string", ...);
```

is exactly equivalent to:

```c
ereport(level, (errmsg_internal("format string", ...)));
```

Notice that the SQLSTATE error code is always defaulted, and the message string is not subject to translation. Therefore, `elog` should be used only for internal errors and low-level debug logging. Any message that is likely to be of interest to ordinary users should go through `ereport`. Nonetheless, there are enough internal “cannot happen” error checks in the system that `elog` is still widely used; it is preferred for those messages for its notational simplicity.

Advice about writing good error messages can be found in Section 52.3.

52.3. Error Message Style Guide

This style guide is offered in the hope of maintaining a consistent, user-friendly style throughout all the messages generated by PostgreSQL.
52.3.1. What Goes Where

The primary message should be short, factual, and avoid reference to implementation details such as specific function names. “Short” means “should fit on one line under normal conditions”. Use a detail message if needed to keep the primary message short, or if you feel a need to mention implementation details such as the particular system call that failed. Both primary and detail messages should be factual. Use a hint message for suggestions about what to do to fix the problem, especially if the suggestion might not always be applicable.

For example, instead of:

```
IpcMemoryCreate: shmget(key=%d, size=%u, 0%o) failed: %m
```

(plus a long addendum that is basically a hint)

write:

```
Primary: could not create shared memory segment: %m
Detail: Failed syscall was shmget(key=%d, size=%u, 0%o).
Hint: the addendum
```

Rationale: keeping the primary message short helps keep it to the point, and lets clients lay out screen space on the assumption that one line is enough for error messages. Detail and hint messages can be relegated to a verbose mode, or perhaps a pop-up error-details window. Also, details and hints would normally be suppressed from the server log to save space. Reference to implementation details is best avoided since users aren’t expected to know the details.

52.3.2. Formatting

Don’t put any specific assumptions about formatting into the message texts. Expect clients and the server log to wrap lines to fit their own needs. In long messages, newline characters (`\n`) can be used to indicate suggested paragraph breaks. Don’t end a message with a newline. Don’t use tabs or other formatting characters. (In error context displays, newlines are automatically added to separate levels of context such as function calls.)

Rationale: Messages are not necessarily displayed on terminal-type displays. In GUI displays or browsers these formatting instructions are at best ignored.

52.3.3. Quotation Marks

English text should use double quotes when quoting is appropriate. Text in other languages should consistently use one kind of quotes that is consistent with publishing customs and computer output of other programs.

Rationale: The choice of double quotes over single quotes is somewhat arbitrary, but tends to be the preferred use. Some have suggested choosing the kind of quotes depending on the type of object according to SQL conventions (namely, strings single quoted, identifiers double quoted). But this is a language-internal technical issue that many users aren’t even familiar with, it won’t scale to other kinds of quoted terms, it doesn’t translate to other languages, and it’s pretty pointless, too.
52.3.4. Use of Quotes

Use quotes always to delimit file names, user-supplied identifiers, and other variables that might contain words. Do not use them to mark up variables that will not contain words (for example, operator names).

There are functions in the backend that will double-quote their own output at need (for example, format_type_be()). Do not put additional quotes around the output of such functions.

Rationale: Objects can have names that create ambiguity when embedded in a message. Be consistent about denoting where a plugged-in name starts and ends. But don’t clutter messages with unnecessary or duplicate quote marks.

52.3.5. Grammar and Punctuation

The rules are different for primary error messages and for detail/hint messages:

Primary error messages: Do not capitalize the first letter. Do not end a message with a period. Do not even think about ending a message with an exclamation point.

Detail and hint messages: Use complete sentences, and end each with a period. Capitalize the first word of sentences. Put two spaces after the period if another sentence follows (for English text; might be inappropriate in other languages).

Error context strings: Do not capitalize the first letter and do not end the string with a period. Context strings should normally not be complete sentences.

Rationale: Avoiding punctuation makes it easier for client applications to embed the message into a variety of grammatical contexts. Often, primary messages are not grammatically complete sentences anyway. (And if they’re long enough to be more than one sentence, they should be split into primary and detail parts.) However, detail and hint messages are longer and might need to include multiple sentences. For consistency, they should follow complete-sentence style even when there’s only one sentence.

52.3.6. Upper Case vs. Lower Case

Use lower case for message wording, including the first letter of a primary error message. Use upper case for SQL commands and key words if they appear in the message.

Rationale: It’s easier to make everything look more consistent this way, since some messages are complete sentences and some not.

52.3.7. Avoid Passive Voice

Use the active voice. Use complete sentences when there is an acting subject (“A could not do B”). Use telegram style without subject if the subject would be the program itself; do not use “I” for the program.

Rationale: The program is not human. Don’t pretend otherwise.
52.3.8. Present vs. Past Tense

Use past tense if an attempt to do something failed, but could perhaps succeed next time (perhaps after fixing some problem). Use present tense if the failure is certainly permanent.

There is a nontrivial semantic difference between sentences of the form:

could not open file "%s": %m

and:

cannot open file "%s"

The first one means that the attempt to open the file failed. The message should give a reason, such as “disk full” or “file doesn’t exist”. The past tense is appropriate because next time the disk might not be full anymore or the file in question might exist.

The second form indicates that the functionality of opening the named file does not exist at all in the program, or that it’s conceptually impossible. The present tense is appropriate because the condition will persist indefinitely.

Rationale: Granted, the average user will not be able to draw great conclusions merely from the tense of the message, but since the language provides us with a grammar we should use it correctly.

52.3.9. Type of the Object

When citing the name of an object, state what kind of object it is.

Rationale: Otherwise no one will know what “foo.bar.baz” refers to.

52.3.10. Brackets

Square brackets are only to be used (1) in command synopses to denote optional arguments, or (2) to denote an array subscript.

Rationale: Anything else does not correspond to widely-known customary usage and will confuse people.

52.3.11. Assembling Error Messages

When a message includes text that is generated elsewhere, embed it in this style:

could not open file %s: %m

Rationale: It would be difficult to account for all possible error codes to paste this into a single smooth sentence, so some sort of punctuation is needed. Putting the embedded text in parentheses has also been suggested, but it’s unnatural if the embedded text is likely to be the most important part of the message, as is often the case.
52.3.12. Reasons for Errors

Messages should always state the reason why an error occurred. For example:

BAD: could not open file %s
BETTER: could not open file %s (I/O failure)

If no reason is known you better fix the code.

52.3.13. Function Names

Don’t include the name of the reporting routine in the error text. We have other mechanisms for finding that out when needed, and for most users it’s not helpful information. If the error text doesn’t make as much sense without the function name, reword it.

BAD: pg_atoi: error in "z": cannot parse "z"
BETTER: invalid input syntax for integer: "z"

Avoid mentioning called function names, either; instead say what the code was trying to do:

BAD: open() failed: %m
BETTER: could not open file %s: %m

If it really seems necessary, mention the system call in the detail message. (In some cases, providing the actual values passed to the system call might be appropriate information for the detail message.)

Rationale: Users don’t know what all those functions do.

52.3.14. Tricky Words to Avoid

Unable. “Unable” is nearly the passive voice. Better use “cannot” or “could not”, as appropriate.

Bad. Error messages like “bad result” are really hard to interpret intelligently. It’s better to write why the result is “bad”, e.g., “invalid format”.

Illegal. “Illegal” stands for a violation of the law, the rest is “invalid”. Better yet, say why it’s invalid.

Unknown. Try to avoid “unknown”. Consider “error: unknown response”. If you don’t know what the response is, how do you know it’s erroneous? “Unrecognized” is often a better choice. Also, be sure to include the value being complained of.

BAD: unknown node type
BETTER: unrecognized node type: 42

Find vs. Exists. If the program uses a nontrivial algorithm to locate a resource (e.g., a path search) and that algorithm fails, it is fair to say that the program couldn’t “find” the resource. If, on the other hand, the expected location of the resource is known but the program cannot access it there then say that the resource doesn’t “exist”. Using “find” in this case sounds weak and confuses the issue.

May vs. Can vs. Might. “May” suggests permission (e.g., “You may borrow my rake.”), and has little use in documentation or error messages. “Can” suggests ability (e.g., “I can lift that log.”), and “might” suggests possibility (e.g., “It might rain today.”). Using the proper word clarifies meaning and assists translation.

Contractions. Avoid contractions, like “can’t”; use “cannot” instead.
52.3.15. Proper Spelling

Spell out words in full. For instance, avoid:

- spec
- stats
- parens
- auth
- xact

Rationale: This will improve consistency.

52.3.16. Localization

Keep in mind that error message texts need to be translated into other languages. Follow the guidelines in Section 53.2.2 to avoid making life difficult for translators.

52.4. Miscellaneous Coding Conventions

52.4.1. C Standard

Code in PostgreSQL should only rely on language features available in the C89 standard. That means a conforming C89 compiler has to be able to compile postgres, at least aside from a few platform dependent pieces. Features from later revision of the C standard or compiler specific features can be used, if a fallback is provided.

For example static inline and _Static_assert() are currently used, even though they are from newer revisions of the C standard. If not available we respectively fall back to defining the functions without inline, and to using a C89 compatible replacement that performs the same checks, but emits rather cryptic messages.

52.4.2. Function-Like Macros and Inline Functions

Both, macros with arguments and static inline functions, may be used. The latter are preferable if there are multiple-evaluation hazards when written as a macro, as e.g. the case with

```c
#define Max(x, y) ((x) > (y) ? (x) : (y))
```

or when the macro would be very long. In other cases it’s only possible to use macros, or at least easier. For example because expressions of various types need to be passed to the macro.

When the definition of an inline function references symbols (i.e. variables, functions) that are only available as part of the backend, the function may not be visible when included from frontend code.

```c
#ifndef FRONTEND
static inline MemoryContext
```

2092
MemoryContextSwitchTo(MemoryContext context)
{
    MemoryContext old = CurrentMemoryContext;

    CurrentMemoryContext = context;
    return old;
}
#endif /* FRONTEND */

In this example CurrentMemoryContext, which is only available in the backend, is referenced and
the function thus hidden with a #ifndef FRONTEND. This rule exists because some compilers emit
references to symbols contained in inline functions even if the function is not used.

52.4.3. Writing Signal Handlers

To be suitable to run inside a signal handler code has to be written very carefully. The fundamental
problem is that, unless blocked, a signal handler can interrupt code at any time. If code inside the
signal handler uses the same state as code outside chaos may ensue. As an example consider what
happens if a signal handler tries to acquire a lock that’s already held in the interrupted code.

Barring special arrangements code in signal handlers may only call async-signal safe functions (as
declared in POSIX) and access variables of type volatile sig_atomic_t. A few functions in
postgres are also deemed signal safe, importantly SetLatch().

In most cases signal handlers should do nothing more than note that a signal has arrived, and wake up
code running outside of the handler using a latch. An example of such a handler is the following:

static void
handle_sighup(SIGNAL_ARGS)
{
    int save_errno = errno;

    got_SIGHUP = true;
    SetLatch(MyLatch);

    errno = save_errno;
}

errno is saved and restored because SetLatch() might change it. If that were not done interrupted
code that’s currently inspecting errno might see the wrong value.
Chapter 53. Native Language Support

53.1. For the Translator

PostgreSQL programs (server and client) can issue their messages in your favorite language — if the messages have been translated. Creating and maintaining translated message sets needs the help of people who speak their own language well and want to contribute to the PostgreSQL effort. You do not have to be a programmer at all to do this. This section explains how to help.

53.1.1. Requirements

We won’t judge your language skills — this section is about software tools. Theoretically, you only need a text editor. But this is only in the unlikely event that you do not want to try out your translated messages. When you configure your source tree, be sure to use the `--enable-nls` option. This will also check for the libintl library and the `msgfmt` program, which all end users will need anyway. To try out your work, follow the applicable portions of the installation instructions.

If you want to start a new translation effort or want to do a message catalog merge (described later), you will need the programs `xgettext` and `msgmerge`, respectively, in a GNU-compatible implementation. Later, we will try to arrange it so that if you use a packaged source distribution, you won’t need `xgettext`. (If working from Git, you will still need it.) GNU Gettext 0.10.36 or later is currently recommended.

Your local gettext implementation should come with its own documentation. Some of that is probably duplicated in what follows, but for additional details you should look there.

53.1.2. Concepts

The pairs of original (English) messages and their (possibly) translated equivalents are kept in message catalogs, one for each program (although related programs can share a message catalog) and for each target language. There are two file formats for message catalogs: The first is the “PO” file (for Portable Object), which is a plain text file with special syntax that translators edit. The second is the “MO” file (for Machine Object), which is a binary file generated from the respective PO file and is used while the internationalized program is run. Translators do not deal with MO files; in fact hardly anyone does.

The extension of the message catalog file is to no surprise either `.po` or `.mo`. The base name is either the name of the program it accompanies, or the language the file is for, depending on the situation. This is a bit confusing. Examples are `psql.po` (PO file for psql) or `fr.mo` (MO file in French).

The file format of the PO files is illustrated here:

```plaintext
# comment
msgid "original string"
msgstr "translated string"
msgid "more original"
```
Chapter 53. Native Language Support

msgstr "another translated"
"string can be broken up like this"

... 

The msgid’s are extracted from the program source. (They need not be, but this is the most common
way.) The msgstr lines are initially empty and are filled in with useful strings by the translator. The
strings can contain C-style escape characters and can be continued across lines as illustrated. (The
next line must start at the beginning of the line.)

The # character introduces a comment. If whitespace immediately follows the # character, then this
is a comment maintained by the translator. There can also be automatic comments, which have a
non-whitespace character immediately following the #. These are maintained by the various tools that
operate on the PO files and are intended to aid the translator.

#. automatic comment
#: filename.c:1023
#, flags, flags

The #. style comments are extracted from the source file where the message is used. Possibly the
programmer has inserted information for the translator, such as about expected alignment. The #:
comment indicates the exact location(s) where the message is used in the source. The translator need
not look at the program source, but can if there is doubt about the correct translation. The #, comments
contain flags that describe the message in some way. There are currently two flags: fuzzy is set if the
message has possibly been outdated because of changes in the program source. The translator can then
verify this and possibly remove the fuzzy flag. Note that fuzzy messages are not made available to
the end user. The other flag is c-format, which indicates that the message is a printf-style format
template. This means that the translation should also be a format string with the same number and
type of placeholders. There are tools that can verify this, which key off the c-format flag.

53.1.3. Creating and Maintaining Message Catalogs

OK, so how does one create a “blank” message catalog? First, go into the directory that contains the
program whose messages you want to translate. If there is a file nls.mk, then this program has been
prepared for translation.

If there are already some .po files, then someone has already done some translation work. The files are
named language.po, where language is the ISO 639-1 two-letter language code (in lower case),
e.g., fr.po for French. If there is really a need for more than one translation effort per language
then the files can also be named language_region.po where region is the ISO 3166-1 two-letter
country code (in upper case), e.g., pt_BR.po for Portuguese in Brazil. If you find the language you
wanted you can just start working on that file.

If you need to start a new translation effort, then first run the command:

make init-po

This will create a file programe.pot (.pot to distinguish it from PO files that are “in production”.
The T stands for “template”.) Copy this file to language.po and edit it. To make it known that the
new language is available, also edit the file nls.mk and add the language (or language and country)
code to the line that looks like:

2. http://www.iso.org/iso/country_names_and_code_elements

2095
AVAIL_LANGUAGES := de fr

(Other languages can appear, of course.)

As the underlying program or library changes, messages might be changed or added by the program-
mers. In this case you do not need to start from scratch. Instead, run the command:

\texttt{make update-po}

which will create a new blank message catalog file (the pot file you started with) and will merge it
with the existing PO files. If the merge algorithm is not sure about a particular message it marks it
“fuzzy” as explained above. The new PO file is saved with a .po.new extension.

53.1.4. Editing the PO Files

The PO files can be edited with a regular text editor. The translator should only change the area
between the quotes after the \texttt{msgstr} directive, add comments, and alter the fuzzy flag. There is (un-
surprisingly) a PO mode for Emacs, which I find quite useful.

The PO files need not be completely filled in. The software will automatically fall back to the original
string if no translation (or an empty translation) is available. It is no problem to submit incomplete
translations for inclusions in the source tree; that gives room for other people to pick up your work.
However, you are encouraged to give priority to removing fuzzy entries after doing a merge. Remem-
ber that fuzzy entries will not be installed; they only serve as reference for what might be the right
translation.

Here are some things to keep in mind while editing the translations:

- Make sure that if the original ends with a newline, the translation does, too. Similarly for tabs, etc.
- If the original is a \texttt{printf} format string, the translation also needs to be. The translation also needs
to have the same format specifiers in the same order. Sometimes the natural rules of the language
make this impossible or at least awkward. In that case you can modify the format specifiers like
this:

\begin{verbatim}
msgstr "Die Datei %2$s hat %1$u Zeichen."
\end{verbatim}

Then the first placeholder will actually use the second argument from the list. The \textit{digits} needs
to follow the % immediately, before any other format manipulators. (This feature really exists in
the \texttt{printf} family of functions. You might not have heard of it before because there is little use
for it outside of message internationalization.)

- If the original string contains a linguistic mistake, report that (or fix it yourself in the program
source) and translate normally. The corrected string can be merged in when the program sources
have been updated. If the original string contains a factual mistake, report that (or fix it yourself)
and do not translate it. Instead, you can mark the string with a comment in the PO file.

- Maintain the style and tone of the original string. Specifically, messages that are not sentences
\texttt{(cannot open file %s)} should probably not start with a capital letter (if your language distin-
guishes letter case) or end with a period (if your language uses punctuation marks). It might help
to read Section 52.3.

- If you don’t know what a message means, or if it is ambiguous, ask on the developers’ mailing list.
Chances are that English speaking end users might also not understand it or find it ambiguous, so
it’s best to improve the message.
53.2. For the Programmer

53.2.1. Mechanics

This section describes how to implement native language support in a program or library that is part of the PostgreSQL distribution. Currently, it only applies to C programs.

Adding NLS Support to a Program

1. Insert this code into the start-up sequence of the program:

   ```c
   #ifdef ENABLE_NLS
   #include <locale.h>
   #endif
   ...
   #ifdef ENABLE_NLS
   setlocale(LC_ALL, """);
   bindtextdomain("progname", LOCALEDIR);
   textdomain("progname");
   #endif
   (The progname can actually be chosen freely.)
   
   2. Wherever a message that is a candidate for translation is found, a call to gettext() needs to be inserted. E.g.:

   ```c
   fprintf(stderr, "panic level %d\n", lvl);
   ```

   would be changed to:

   ```c
   fprintf(stderr, gettext("panic level %d\n"), lvl);
   ```

   (gettext is defined as a no-op if NLS support is not configured.)

   This tends to add a lot of clutter. One common shortcut is to use:

   ```c
   #define _(x) gettext(x)
   ```

   Another solution is feasible if the program does much of its communication through one or a few functions, such as ereport() in the backend. Then you make this function call gettext internally on all input strings.

   3. Add a file nls.mk in the directory with the program sources. This file will be read as a makefile. The following variable assignments need to be made here:

   CATALOG_NAME
   The program name, as provided in the textdomain() call.

   AVAIL_LANGUAGES
   List of provided translations — initially empty.

   GETTEXT_FILES
   List of files that contain translatable strings, i.e., those marked with gettext or an alternative solution. Eventually, this will include nearly all source files of the program. If this list gets too long you can make the first “file” be a + and the second word be a file that contains one file name per line.
GETTEXT_TRIGGERS

The tools that generate message catalogs for the translators to work on need to know what
function calls contain translatable strings. By default, only gettext() calls are known. If
you used _ or other identifiers you need to list them here. If the translatable string is not
the first argument, the item needs to be of the form func:2 (for the second argument). If
you have a function that supports pluralized messages, the item should look like func:1,2
(identifying the singular and plural message arguments).

The build system will automatically take care of building and installing the message catalogs.

53.2.2. Message-writing Guidelines

Here are some guidelines for writing messages that are easily translatable.

- Do not construct sentences at run-time, like:

  printf("Files were %s.\n", flag ? "copied" : "removed");

  The word order within the sentence might be different in other languages. Also, even if you re-
  member to call gettext() on each fragment, the fragments might not translate well separately.
  It’s better to duplicate a little code so that each message to be translated is a coherent whole. Only
  numbers, file names, and such-like run-time variables should be inserted at run time into a message
text.

- For similar reasons, this won’t work:

  printf("copied %d file%s", n, n!=1 ? "s" : "");

  because it assumes how the plural is formed. If you figured you could solve it like this:

  if (n==1)
      printf("copied 1 file");
  else
      printf("copied %d files", n):

  then be disappointed. Some languages have more than two forms, with some peculiar rules. It’s
  often best to design the message to avoid the issue altogether, for instance like this:

  printf("number of copied files: %d", n);

  If you really want to construct a properly pluralized message, there is support for this, but it’s a
  bit awkward. When generating a primary or detail error message in ereport(), you can write
  something like this:

  errmsg_plural("copied %d file",
         "copied %d files",
         n,
         n)

  The first argument is the format string appropriate for English singular form, the second is the
  format string appropriate for English plural form, and the third is the integer control value that
determines which plural form to use. Subsequent arguments are formatted per the format string as
usual. (Normally, the pluralization control value will also be one of the values to be formatted, so
it has to be written twice.) In English it only matters whether n is 1 or not 1, but in other languages
there can be many different plural forms. The translator sees the two English forms as a group and
has the opportunity to supply multiple substitute strings, with the appropriate one being selected
based on the run-time value of n.
If you need to pluralize a message that isn’t going directly to an `errmsg` or `errdetail` report, you have to use the underlying function `ngettext`. See the gettext documentation.

- If you want to communicate something to the translator, such as about how a message is intended to line up with other output, precede the occurrence of the string with a comment that starts with `translator`, e.g.:

  ```c
  /* translator: This message is not what it seems to be. */
  ```

  These comments are copied to the message catalog files so that the translators can see them.
Chapter 54. Writing A Procedural Language Handler

All calls to functions that are written in a language other than the current “version 1” interface for compiled languages (this includes functions in user-defined procedural languages, functions written in SQL, and functions using the version 0 compiled language interface) go through a call handler function for the specific language. It is the responsibility of the call handler to execute the function in a meaningful way, such as by interpreting the supplied source text. This chapter outlines how a new procedural language’s call handler can be written.

The call handler for a procedural language is a “normal” function that must be written in a compiled language such as C, using the version-1 interface, and registered with PostgreSQL as taking no arguments and returning the type language_handler. This special pseudotype identifies the function as a call handler and prevents it from being called directly in SQL commands. For more details on C language calling conventions and dynamic loading, see Section 36.9.

The call handler is called in the same way as any other function: It receives a pointer to a FunctionCallInfoData struct containing argument values and information about the called function, and it is expected to return a Datum result (and possibly set the isnull field of the FunctionCallInfoData structure, if it wishes to return an SQL null result). The difference between a call handler and an ordinary callee function is that the flinfo->fn_oid field of the FunctionCallInfoData structure will contain the OID of the actual function to be called, not of the call handler itself. The call handler must use this field to determine which function to execute. Also, the passed argument list has been set up according to the declaration of the target function, not of the call handler.

It’s up to the call handler to fetch the entry of the function from the pg_proc system catalog and to analyze the argument and return types of the called function. The AS clause from the CREATE FUNCTION command for the function will be found in the prosrc column of the pg_proc row. This is commonly source text in the procedural language, but in theory it could be something else, such as a path name to a file, or anything else that tells the call handler what to do in detail.

Often, the same function is called many times per SQL statement. A call handler can avoid repeated lookups of information about the called function by using the flinfo->fn_extra field. This will initially be NULL, but can be set by the call handler to point at information about the called function. On subsequent calls, if flinfo->fn_extra is already non-NULL then it can be used and the information lookup step skipped. The call handler must make sure that flinfo->fn_extra is made to point at memory that will live at least until the end of the current query, since an FmgrInfo data structure could be kept that long. One way to do this is to allocate the extra data in the memory context specified by flinfo->fn_mcxt; such data will normally have the same lifespan as the FmgrInfo itself. But the handler could also choose to use a longer-lived memory context so that it can cache function definition information across queries.

When a procedural-language function is invoked as a trigger, no arguments are passed in the usual way, but the FunctionCallInfoData’s context field points at a TriggerData structure, rather than being NULL as it is in a plain function call. A language handler should provide mechanisms for procedural-language functions to get at the trigger information.

This is a template for a procedural-language handler written in C:

#include "postgres.h"
#include "executor/spi.h"
#include "commands/trigger.h"
#include "fmgr.h"
#include "access/heapam.h"
#include "utils/syscache.h"
#include "catalog/pg_proc.h"
#include "catalog/pg_type.h"

#ifdef PG_MODULE_MAGIC
    PG_MODULE_MAGIC;
#endif

PG_FUNCTION_INFO_V1(plsample_call_handler);

Datum plsample_call_handler(PG_FUNCTION_ARGS)
{
    Datum retval;

    if (CALLED_AS_TRIGGER(fcinfo))
    {
        /*
        * Called as a trigger procedure
        */
        TriggerData *trigdata = (TriggerData *) fcinfo->context;
        retval = ...
    }
    else
    {
        /*
        * Called as a function
        */
        retval = ...
    }

    return retval;
}

Only a few thousand lines of code have to be added instead of the dots to complete the call handler.

After having compiled the handler function into a loadable module (see Section 36.9.6), the following
commands then register the sample procedural language:

CREATE FUNCTION plsample_call_handler() RETURNS language_handler
    AS 'filename'
    LANGUAGE C;
CREATE LANGUAGE plsample
    HANDLER plsample_call_handler;

Although providing a call handler is sufficient to create a minimal procedural language, there are two
other functions that can optionally be provided to make the language more convenient to use. These
are a validator and an inline handler. A validator can be provided to allow language-specific checking
to be done during CREATE FUNCTION. An inline handler can be provided to allow the language to
support anonymous code blocks executed via the DO command.
If a validator is provided by a procedural language, it must be declared as a function taking a single parameter of type oid. The validator’s result is ignored, so it is customarily declared to return void. The validator will be called at the end of a CREATE FUNCTION command that has created or updated a function written in the procedural language. The passed-in OID is the OID of the function’s pg_proc row. The validator must fetch this row in the usual way, and do whatever checking is appropriate. First, call CheckFunctionValidatorAccess() to diagnose explicit calls to the validator that the user could not achieve through CREATE FUNCTION. Typical checks then include verifying that the function’s argument and result types are supported by the language, and that the function’s body is syntactically correct in the language. If the validator finds the function to be okay, it should just return. If it finds an error, it should report that via the normal ereport() error reporting mechanism. Throwing an error will force a transaction rollback and thus prevent the incorrect function definition from being committed.

Validator functions should typically honor the check_function_bodies parameter: if it is turned off then any expensive or context-sensitive checking should be skipped. If the language provides for code execution at compilation time, the validator must suppress checks that would induce such execution. In particular, this parameter is turned off by pg_dump so that it can load procedural language functions without worrying about side effects or dependencies of the function bodies on other database objects. (Because of this requirement, the call handler should avoid assuming that the validator has fully checked the function. The point of having a validator is not to let the call handler omit checks, but to notify the user immediately if there are obvious errors in a CREATE FUNCTION command.) While the choice of exactly what to check is mostly left to the discretion of the validator function, note that the core CREATE FUNCTION code only executes SET clauses attached to a function when check_function_bodies is on. Therefore, checks whose results might be affected by GUC parameters definitely should be skipped when check_function_bodies is off, to avoid false failures when reloading a dump.

If an inline handler is provided by a procedural language, it must be declared as a function taking a single parameter of type internal. The inline handler’s result is ignored, so it is customarily declared to return void. The inline handler will be called when a DO statement is executed specifying the procedural language. The parameter actually passed is a pointer to an InlineCodeBlock struct, which contains information about the DO statement’s parameters, in particular the text of the anonymous code block to be executed. The inline handler should execute this code and return.

It’s recommended that you wrap all these function declarations, as well as the CREATE LANGUAGE command itself, into an extension so that a simple CREATE EXTENSION command is sufficient to install the language. See Section 36.15 for information about writing extensions.

The procedural languages included in the standard distribution are good references when trying to write your own language handler. Look into the src/pl subdirectory of the source tree. The CREATE LANGUAGE reference page also has some useful details.
Chapter 55. Writing A Foreign Data Wrapper

All operations on a foreign table are handled through its foreign data wrapper, which consists of a set of functions that the core server calls. The foreign data wrapper is responsible for fetching data from the remote data source and returning it to the PostgreSQL executor. If updating foreign tables is to be supported, the wrapper must handle that, too. This chapter outlines how to write a new foreign data wrapper.

The foreign data wrappers included in the standard distribution are good references when trying to write your own. Look into the contrib subdirectory of the source tree. The CREATE FOREIGN DATA WRAPPER reference page also has some useful details.

Note: The SQL standard specifies an interface for writing foreign data wrappers. However, PostgreSQL does not implement that API, because the effort to accommodate it into PostgreSQL would be large, and the standard API hasn’t gained wide adoption anyway.

55.1. Foreign Data Wrapper Functions

The FDW author needs to implement a handler function, and optionally a validator function. Both functions must be written in a compiled language such as C, using the version-1 interface. For details on C language calling conventions and dynamic loading, see Section 36.9.

The handler function simply returns a struct of function pointers to callback functions that will be called by the planner, executor, and various maintenance commands. Most of the effort in writing an FDW is in implementing these callback functions. The handler function must be registered with PostgreSQL as taking no arguments and returning the special pseudo-type `fdw_handler`. The callback functions are plain C functions and are not visible or callable at the SQL level. The callback functions are described in Section 55.2.

The validator function is responsible for validating options given in CREATE and ALTER commands for its foreign data wrapper, as well as foreign servers, user mappings, and foreign tables using the wrapper. The validator function must be registered as taking two arguments, a text array containing the options to be validated, and an OID representing the type of object the options are associated with (in the form of the OID of the system catalog the object would be stored in, either `ForeignDataWrapperRelationId`, `ForeignServerRelationId`, `UserMappingRelationId`, or `ForeignTableRelationId`). If no validator function is supplied, options are not checked at object creation time or object alteration time.

55.2. Foreign Data Wrapper Callback Routines

The FDW handler function returns a palloc’d `FdwRoutine` struct containing pointers to the callback functions described below. The scan-related functions are required, the rest are optional.

The `FdwRoutine` struct type is declared in `src/include/foreign/fdwapi.h`, which see for additional details.
Chapter 55. Writing A Foreign Data Wrapper

55.2.1. FDW Routines For Scanning Foreign Tables

```c
void GetForeignRelSize (PlannerInfo *root, 
    RelOptInfo *baserel, 
    Oid foreigntableid);
```

Obtain relation size estimates for a foreign table. This is called at the beginning of planning for a 
query that scans a foreign table. root is the planner’s global information about the query; baserel 
is the planner’s information about this table; and foreigntableid is the pg_class OID of the 
foreign table. (foreigntableid could be obtained from the planner data structures, but it’s passed 
explicitly to save effort.)

This function should update baserel->rows to be the expected number of rows returned by 
the table scan, after accounting for the filtering done by the restriction quals. The initial value of 
baserel->rows is just a constant default estimate, which should be replaced if at all possible. The 
function may also choose to update baserel->width if it can compute a better estimate of the 
average result row width.

See Section 55.4 for additional information.

```c
void GetForeignPaths (PlannerInfo *root, 
    RelOptInfo *baserel, 
    Oid foreigntableid);
```

Create possible access paths for a scan on a foreign table. This is called during query planning. The 
parameters are the same as for GetForeignRelSize, which has already been called.

This function must generate at least one access path (ForeignPath node) for a scan on the foreign 
table and must call add_path to add each such path to baserel->pathlist. It’s recommended 
to use create_foreignscan_path to build the ForeignPath nodes. The function can generate 
multiple access paths, e.g., a path which has valid pathkeys to represent a pre-sorted result. Each 
access path must contain cost estimates, and can contain any FDW-private information that is needed 
to identify the specific scan method intended.

See Section 55.4 for additional information.

```c
ForeignScan * GetForeignPlan (PlannerInfo *root, 
    RelOptInfo *baserel, 
    Oid foreigntableid, 
    ForeignPath *best_path, 
    List *tlist, 
    List *scan_clauses, 
    Plan *outer_plan);
```

Create a ForeignScan plan node from the selected foreign access path. This is called at the 
end of query planning. The parameters are as for GetForeignRelSize, plus the selected 
ForeignPath (previously produced by GetForeignPaths, GetForeignJoinPaths, or 
GetForeignUpperPaths), the target list to be emitted by the plan node, the restriction clauses 
to be enforced by the plan node, and the outer subplan of the ForeignScan, which is used for 
rechecks performed by RecheckForeignScan. (If the path is for a join rather than a base relation, 
foreigntableid is InvalidOid.)
This function must create and return a ForeignScan plan node; it’s recommended to use make_foreignscan to build the ForeignScan node.

See Section 55.4 for additional information.

```c
void BeginForeignScan (ForeignScanState *node, int eflags);
```

Begin executing a foreign scan. This is called during executor startup. It should perform any initialization needed before the scan can start, but not start executing the actual scan (that should be done upon the first call to IterateForeignScan). The ForeignScanState node has already been created, but its fdw_state field is still NULL. Information about the table to scan is accessible through the ForeignScanState node (in particular, from the underlying ForeignScan plan node, which contains any FDW-private information provided by GetForeignPlan). eflags contains flag bits describing the executor’s operating mode for this plan node.

Note that when (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, this function should not perform any externally-visible actions; it should only do the minimum required to make the node state valid for ExplainForeignScan and EndForeignScan.

```c
TupleTableSlot * IterateForeignScan (ForeignScanState *node);
```

Fetch one row from the foreign source, returning it in a tuple table slot (the node’s ScanTupleSlot should be used for this purpose). Return NULL if no more rows are available. The tuple table slot infrastructure allows either a physical or virtual tuple to be returned; in most cases the latter choice is preferable from a performance standpoint. Note that this is called in a short-lived memory context that will be reset between invocations. Create a memory context in BeginForeignScan if you need longer-lived storage, or use the es_query_cxt of the node’s EState.

The rows returned must match the fdw_scan_tlist target list if one was supplied, otherwise they must match the row type of the foreign table being scanned. If you choose to optimize away fetching columns that are not needed, you should insert nulls in those column positions, or else generate a fdw_scan_tlist list with those columns omitted.

Note that PostgreSQL’s executor doesn’t care whether the rows returned violate any constraints that were defined on the foreign table — but the planner does care, and may optimize queries incorrectly if there are rows visible in the foreign table that do not satisfy a declared constraint. If a constraint is violated when the user has declared that the constraint should hold true, it may be appropriate to raise an error (just as you would need to do in the case of a data type mismatch).

```c
void ReScanForeignScan (ForeignScanState *node);
```

Restart the scan from the beginning. Note that any parameters the scan depends on may have changed value, so the new scan does not necessarily return exactly the same rows.

```c
void EndForeignScan (ForeignScanState *node);
```

End the scan and release resources. It is normally not important to release palloc’d memory, but for example open files and connections to remote servers should be cleaned up.
55.2.2. FDW Routines For Scanning Foreign Joins

If an FDW supports performing foreign joins remotely (rather than by fetching both tables’ data and doing the join locally), it should provide this callback function:

```c
void
GetForeignJoinPaths (PlannerInfo *root,
             RelOptInfo *joinrel,
             RelOptInfo *outerrel,
             RelOptInfo *innerrel,
             JoinType jointype,
             JoinPathExtraData *extra);
```

Create possible access paths for a join of two (or more) foreign tables that all belong to the same foreign server. This optional function is called during query planning. As with `GetForeignPaths`, this function should generate `ForeignPath` path(s) for the supplied `joinrel`, and call `add_path` to add these paths to the set of paths considered for the join. But unlike `GetForeignPaths`, it is not necessary that this function succeed in creating at least one path, since paths involving local joining are always possible.

Note that this function will be invoked repeatedly for the same join relation, with different combinations of inner and outer relations; it is the responsibility of the FDW to minimize duplicated work.

If a `ForeignPath` path is chosen for the join, it will represent the entire join process; paths generated for the component tables and subsidiary joins will not be used. Subsequent processing of the join path proceeds much as it does for a path scanning a single foreign table. One difference is that the `scanrelid` of the resulting `ForeignScan` plan node should be set to zero, since there is no single relation that it represents; instead, the `fs_relids` field of the `ForeignScan` node represents the set of relations that were joined. (The latter field is set up automatically by the core planner code, and need not be filled by the FDW.) Another difference is that, because the column list for a remote join cannot be found from the system catalogs, the FDW must fill `fdw_scan_tlist` with an appropriate list of `TargetEntry` nodes, representing the set of columns it will supply at run time in the tuples it returns.

See Section 55.4 for additional information.

55.2.3. FDW Routines For Planning Post-Scan/Join Processing

If an FDW supports performing remote post-scan/join processing, such as remote aggregation, it should provide this callback function:

```c
void
GetForeignUpperPaths (PlannerInfo *root,
              UpperRelationKind stage,
              RelOptInfo *input_rel,
              RelOptInfo *output_rel);
```

Create possible access paths for `upper relation` processing, which is the planner’s term for all post-scan/join query processing, such as aggregation, window functions, sorting, and table updates. This optional function is called during query planning. Currently, it is called only if all base relation(s) involved in the query belong to the same FDW. This function should generate `ForeignPath` path(s) for any post-scan/join processing that the FDW knows how to perform remotely, and call `add_path` to add these paths to the indicated upper relation. As with `GetForeignJoinPaths`, it is not necessary
that this function succeed in creating any paths, since paths involving local processing are always possible.

The `stage` parameter identifies which post-scan/join step is currently being considered. `output_rel` is the upper relation that should receive paths representing computation of this step, and `input_rel` is the relation representing the input to this step. (Note that `ForeignPath` paths added to `output_rel` would typically not have any direct dependency on paths of the `input_rel`, since their processing is expected to be done externally. However, examining paths previously generated for the previous processing step can be useful to avoid redundant planning work.)

See Section 55.4 for additional information.

### 55.2.4. FDW Routines For Updating Foreign Tables

If an FDW supports writable foreign tables, it should provide some or all of the following callback functions depending on the needs and capabilities of the FDW:

```c
void AddForeignUpdateTargets (Query *parsetree, RangeTblEntry *target_rte, Relation target_relation);
```

`UPDATE` and `DELETE` operations are performed against rows previously fetched by the table-scanning functions. The FDW may need extra information, such as a row ID or the values of primary-key columns, to ensure that it can identify the exact row to update or delete. To support that, this function can add extra hidden, or “junk”, target columns to the list of columns that are to be retrieved from the foreign table during an `UPDATE` or `DELETE`.

To do that, add `TargetEntry` items to `parsetree->targetList`, containing expressions for the extra values to be fetched. Each such entry must be marked `resjunk = true`, and must have a distinct `resname` that will identify it at execution time. Avoid using names matching `ctidN`, `wholerow`, or `wholerowN`, as the core system can generate junk columns of these names. If the extra expressions are more complex than simple Vars, they must be run through `eval_const_expressions` before adding them to the `targetlist`.

Although this function is called during planning, the information provided is a bit different from that available to other planning routines. `parsetree` is the parse tree for the `UPDATE` or `DELETE` command, while `target_rte` and `target_relation` describe the target foreign table.

If the `AddForeignUpdateTargets` pointer is set to `NULL`, no extra target expressions are added. (This will make it impossible to implement `DELETE` operations, though `UPDATE` may still be feasible if the FDW relies on an unchanging primary key to identify rows.)

```c
List * PlanForeignModify (PlannerInfo *root, ModifyTable *plan, Index resultRelation, int subplan_index);
```

Perform any additional planning actions needed for an insert, update, or delete on a foreign table. This function generates the FDW-private information that will be attached to the `ModifyTable` plan node that performs the update action. This private information must have the form of a `List`, and will be delivered to `BeginForeignModify` during the execution stage.
root is the planner’s global information about the query. plan is the ModifyTable plan node, which is complete except for the fdwPrivLists field. resultRelation identifies the target foreign table by its range table index. subplan_index identifies which target of the ModifyTable plan node this is, counting from zero; use this if you want to index into plan->plans or other substructure of the plan node.

See Section 55.4 for additional information.

If the PlanForeignModify pointer is set to NULL, no additional plan-time actions are taken, and the fdw_private list delivered to BeginForeignModify will be NIL.

```c
void BeginForeignModify (ModifyTableState *mtstate,
            ResultRelInfo *rinfo,
            List *fdw_private,
            int subplan_index,
            int eflags);
```

Begin executing a foreign table modification operation. This routine is called during executor startup. It should perform any initialization needed prior to the actual table modifications. Subsequently, ExecForeignInsert, ExecForeignUpdate or ExecForeignDelete will be called for each tuple to be inserted, updated, or deleted.

mtstate is the overall state of the ModifyTable plan node being executed; global data about the plan and execution state is available via this structure. rinfo is the ResultRelInfo struct describing the target foreign table. (The ri_FdwState field of ResultRelInfo is available for the FDW to store any private state it needs for this operation.) fdw_private contains the private data generated by PlanForeignModify, if any. subplan_index identifies which target of the ModifyTable plan node this is. eflags contains flag bits describing the executor’s operating mode for this plan node.

Note that when (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, this function should not perform any externally-visible actions; it should only do the minimum required to make the node state valid for ExplainForeignModify and EndForeignModify.

If the BeginForeignModify pointer is set to NULL, no action is taken during executor startup.

```c
TupleTableSlot *
ExecForeignInsert (EState *estate,
            ResultRelInfo *rinfo,
            TupleTableSlot *slot,
            TupleTableSlot *planSlot);
```

Insert one tuple into the foreign table. estate is global execution state for the query. rinfo is the ResultRelInfo struct describing the target foreign table. slot contains the tuple to be inserted; it will match the row-type definition of the foreign table. planSlot contains the tuple that was generated by the ModifyTable plan node’s subplan; it differs from slot in possibly containing additional “junk” columns. (The planSlot is typically of little interest for INSERT cases, but is provided for completeness.)

The return value is either a slot containing the data that was actually inserted (this might differ from the data supplied, for example as a result of trigger actions), or NULL if no row was actually inserted (again, typically as a result of triggers). The passed-in slot can be re-used for this purpose.

The data in the returned slot is used only if the INSERT query has a RETURNING clause or the foreign table has an AFTER ROW trigger. Triggers require all columns, but the FDW could choose to optimize away returning some or all columns depending on the contents of the RETURNING clause. Regardless, some slot must be returned to indicate success, or the query’s reported row count will be wrong.
Chapter 55. Writing A Foreign Data Wrapper

If the `ExecForeignInsert` pointer is set to NULL, attempts to insert into the foreign table will fail with an error message.

```c
TupleTableSlot *
ExecForeignUpdate (EState *estate,
     ResultRelInfo *rinfo,
     TupleTableSlot *slot,
     TupleTableSlot *planSlot);
```

Update one tuple in the foreign table. `estate` is global execution state for the query. `rinfo` is the `ResultRelInfo` struct describing the target foreign table. `slot` contains the new data for the tuple; it will match the row-type definition of the foreign table. `planSlot` contains the tuple that was generated by the `ModifyTable` plan node’s subplan; it differs from `slot` in possibly containing additional “junk” columns. In particular, any junk columns that were requested by `AddForeignUpdateTargets` will be available from this slot.

The return value is either a slot containing the row as it was actually updated (this might differ from the data supplied, for example as a result of trigger actions), or NULL if no row was actually updated (again, typically as a result of triggers). The passed-in `slot` can be re-used for this purpose.

The data in the returned slot is used only if the `UPDATE` query has a `RETURNING` clause or the foreign table has an `AFTER ROW` trigger. Triggers require all columns, but the FDW could choose to optimize away returning some or all columns depending on the contents of the `RETURNING` clause. Regardless, some slot must be returned to indicate success, or the query’s reported row count will be wrong.

If the `ExecForeignUpdate` pointer is set to NULL, attempts to update the foreign table will fail with an error message.

```c
TupleTableSlot *
ExecForeignDelete (EState *estate,
     ResultRelInfo *rinfo,
     TupleTableSlot *slot,
     TupleTableSlot *planSlot);
```

Delete one tuple from the foreign table. `estate` is global execution state for the query. `rinfo` is the `ResultRelInfo` struct describing the target foreign table. `slot` contains nothing useful upon call, but can be used to hold the returned tuple. `planSlot` contains the tuple that was generated by the `ModifyTable` plan node’s subplan; in particular, it will carry any junk columns that were requested by `AddForeignUpdateTargets`. The junk column(s) must be used to identify the tuple to be deleted.

The return value is either a slot containing the row that was deleted, or NULL if no row was deleted (typically as a result of triggers). The passed-in `slot` can be used to hold the tuple to be returned.

The data in the returned slot is used only if the `DELETE` query has a `RETURNING` clause or the foreign table has an `AFTER ROW` trigger. Triggers require all columns, but the FDW could choose to optimize away returning some or all columns depending on the contents of the `RETURNING` clause. Regardless, some slot must be returned to indicate success, or the query’s reported row count will be wrong.

If the `ExecForeignDelete` pointer is set to NULL, attempts to delete from the foreign table will fail with an error message.

```c
void
EndForeignModify (EState *estate,
     ResultRelInfo *rinfo);
```
End the table update and release resources. It is normally not important to release palloc’d memory, but for example open files and connections to remote servers should be cleaned up.

If the \texttt{EndForeignModify} pointer is set to \texttt{NULL}, no action is taken during executor shutdown.

\begin{verbatim}
int IsForeignRelUpdatable (Relation rel);
\end{verbatim}

Report which update operations the specified foreign table supports. The return value should be a bit mask of rule event numbers indicating which operations are supported by the foreign table, using the \texttt{CmdType} enumeration; that is, \((1 \ll \text{CMD\_UPDATE}) = 4\) for UPDATE, \((1 \ll \text{CMD\_INSERT}) = 8\) for INSERT, and \((1 \ll \text{CMD\_DELETE}) = 16\) for DELETE.

If the \texttt{IsForeignRelUpdatable} pointer is set to \texttt{NULL}, foreign tables are assumed to be insertable, updatable, or deletable if the FDW provides \texttt{ExecForeignInsert}, \texttt{ExecForeignUpdate}, or \texttt{ExecForeignDelete} respectively. This function is only needed if the FDW supports some tables that are updatable and some that are not. (Even then, it’s permissible to throw an error in the execution routine instead of checking in this function. However, this function is used to determine updatability for display in the \texttt{information\_schema} views.)

Some inserts, updates, and deletes to foreign tables can be optimized by implementing an alternative set of interfaces. The ordinary interfaces for inserts, updates, and deletes fetch rows from the remote server and then modify those rows one at a time. In some cases, this row-by-row approach is necessary, but it can be inefficient. If it is possible for the foreign server to determine which rows should be modified without actually retrieving them, and if there are no local structures which would affect the operation (row-level local triggers or \texttt{WITH CHECK OPTION} constraints from parent views), then it is possible to arrange things so that the entire operation is performed on the remote server. The interfaces described below make this possible.

\begin{verbatim}
bool PlanDirectModify (PlannerInfo *root, MODIFYTable *plan, Index resultRelation, int subplan_index);
\end{verbatim}

Decide whether it is safe to execute a direct modification on the remote server. If so, return \texttt{true} after performing planning actions needed for that. Otherwise, return \texttt{false}. This optional function is called during query planning. If this function succeeds, \texttt{BeginDirectModify}, \texttt{IterateDirectModify} and \texttt{EndDirectModify} will be called at the execution stage, instead. Otherwise, the table modification will be executed using the table-updating functions described above. The parameters are the same as for \texttt{PlanForeignModify}.

To execute the direct modification on the remote server, this function must rewrite the target sub-plan with a \texttt{ForeignScan} plan node that executes the direct modification on the remote server. The \texttt{operation} field of the \texttt{ForeignScan} must be set to the \texttt{CmdType} enumeration appropriately; that is, \texttt{CMD\_UPDATE} for UPDATE, \texttt{CMD\_INSERT} for INSERT, and \texttt{CMD\_DELETE} for DELETE.

See Section 55.4 for additional information.

If the \texttt{PlanDirectModify} pointer is set to \texttt{NULL}, no attempts to execute a direct modification on the remote server are taken.

\begin{verbatim}
void BeginDirectModify (ForeignScanState *node, int eflags);
\end{verbatim}
Chapter 55. Writing A Foreign Data Wrapper

Prepare to execute a direct modification on the remote server. This is called during executor startup. It should perform any initialization needed prior to the direct modification (that should be done upon the first call to IterateDirectModify). The ForeignScanState node has already been created, but its fdw_state field is still NULL. Information about the table to modify is accessible through the ForeignScanState node (in particular, from the underlying ForeignScan plan node, which contains any FDW-private information provided by PlanDirectModify). eflags contains flag bits describing the executor’s operating mode for this plan node.

Note that when (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, this function should not perform any externally-visible actions; it should only do the minimum required to make the node state valid for ExplainDirectModify and EndDirectModify.

If the BeginDirectModify pointer is set to NULL, no attempts to execute a direct modification on the remote server are taken.

TupleTableSlot * IterateDirectModify (ForeignScanState *node);

When the INSERT, UPDATE or DELETE query doesn’t have a RETURNING clause, just return NULL after a direct modification on the remote server. When the query has the clause, fetch one result containing the data needed for the RETURNING calculation, returning it in a tuple table slot (the node’s ScanTupleSlot should be used for this purpose). The data that was actually inserted, updated or deleted must be stored in the es_result_relation_info->ri_projectReturning->pi_exprContext->ecxt_scantuple of the node’s EState. Return NULL if no more rows are available. Note that this is called in a short-lived memory context that will be reset between invocations. Create a memory context in BeginDirectModify if you need longer-lived storage, or use the es_query_ctx of the node’s EState.

The rows returned must match the fdw_scan_tlist target list if one was supplied, otherwise they must match the row type of the foreign table being updated. If you choose to optimize away fetching columns that are not needed for the RETURNING calculation, you should insert nulls in those column positions, or else generate a fdw_scan_tlist list with those columns omitted.

Whether the query has the clause or not, the query’s reported row count must be incremented by the FDW itself. When the query doesn’t have the clause, the FDW must also increment the row count for the ForeignScanState node in the EXPLAIN ANALYZE case.

If the IterateDirectModify pointer is set to NULL, no attempts to execute a direct modification on the remote server are taken.

void EndDirectModify (ForeignScanState *node);

Clean up following a direct modification on the remote server. It is normally not important to release palloc’d memory, but for example open files and connections to the remote server should be cleaned up.

If the EndDirectModify pointer is set to NULL, no attempts to execute a direct modification on the remote server are taken.

55.2.5. FDW Routines For Row Locking

If an FDW wishes to support late row locking (as described in Section 55.5), it must provide the following callback functions:
RowMarkType
GetForeignRowMarkType (RangeTblEntry *rte,
   LockClauseStrength strength);

Report which row-marking option to use for a foreign table. rte is the RangeTblEntry node for
the table and strength describes the lock strength requested by the relevant FOR UPDATE/SHARE
clause, if any. The result must be a member of the RowMarkType enum type.

This function is called during query planning for each foreign table that appears in an UPDATE,
DELETE, or SELECT FOR UPDATE/SHARE query and is not the target of UPDATE or DELETE.

If the GetForeignRowMarkType pointer is set to NULL, the ROW_MARK_COPY option is always used.
(This implies that RefetchForeignRow will never be called, so it need not be provided either.)

See Section 55.5 for more information.

HeapTuple
RefetchForeignRow (EState *estate,
   ExecRowMark *erm,
   Datum rowid,
   bool *updated);

Re-fetch one tuple from the foreign table, after locking it if required. estate is global execution state
for the query, erm is the ExecRowMark struct describing the target foreign table and the row lock type
(if any) to acquire. rowid identifies the tuple to be fetched. updated is an output parameter.

This function should return a malloc'd copy of the fetched tuple, or NULL if the row lock couldn't be
obtained. The row lock type to acquire is defined by erm->markType, which is the value previously
returned by GetForeignRowMarkType. (ROW_MARK_REFERENCE means to just re-fetch the tuple
without acquiring any lock, and ROW_MARK_COPY will never be seen by this routine.)

In addition, *updated should be set to true if what was fetched was an updated version of the
tuple rather than the same version previously obtained. (If the FDW cannot be sure about this, always
returning true is recommended.)

Note that by default, failure to acquire a row lock should result in raising an error; a NULL return is
only appropriate if the SKIP LOCKED option is specified by erm->waitPolicy.

The rowid is the ctid value previously read for the row to be re-fetched. Although the rowid value
is passed as a Datum, it can currently only be a tid. The function API is chosen in hopes that it may
be possible to allow other data types for row IDs in future.

If the RefetchForeignRow pointer is set to NULL, attempts to re-fetch rows will fail with an error
message.

See Section 55.5 for more information.

bool
RecheckForeignScan (ForeignScanState *node, TupleTableSlot *slot);

Recheck that a previously-returned tuple still matches the relevant scan and join qualifiers, and possi-
bly provide a modified version of the tuple. For foreign data wrappers which do not perform join push-
down, it will typically be more convenient to set this to NULL and instead set fdw_recheck_quals
appropriately. When outer joins are pushed down, however, it isn't sufficient to reapply the checks
relevant to all the base tables to the result tuple, even if all needed attributes are present, because
failure to match some qualifier might result in some attributes going to NULL, rather than in no tuple
being returned. RecheckForeignScan can recheck qualifiers and return true if they are still satisfied
and false otherwise, but it can also store a replacement tuple into the supplied slot.
Chapter 55. Writing A Foreign Data Wrapper

To implement join pushdown, a foreign data wrapper will typically construct an alternative local join plan which is used only for rechecks; this will become the outer subplan of the ForeignScan. When a recheck is required, this subplan can be executed and the resulting tuple can be stored in the slot. This plan need not be efficient since no base table will return more than one row; for example, it may implement all joins as nested loops. The function GetExistingLocalJoinPath may be used to search existing paths for a suitable local join path, which can be used as the alternative local join plan. GetExistingLocalJoinPath searches for an unparameterized path in the path list of the specified join relation. (If it does not find such a path, it returns NULL, in which case a foreign data wrapper may build the local path by itself or may choose not to create access paths for that join.)

55.2.6. FDW Routines for EXPLAIN

void
ExplainForeignScan (ForeignScanState *node,
                     ExplainState *es);

Print additional EXPLAIN output for a foreign table scan. This function can call ExplainPropertyText and related functions to add fields to the EXPLAIN output. The flag fields in es can be used to determine what to print, and the state of the ForeignScanState node can be inspected to provide run-time statistics in the EXPLAIN ANALYZE case.

If the ExplainForeignScan pointer is set to NULL, no additional information is printed during EXPLAIN.

void
ExplainForeignModify (ModifyTableState *mtstate,
                       ResultRelInfo *rinfo,
                       List *fdw_private,
                       int subplan_index,
                       struct ExplainState *es);

Print additional EXPLAIN output for a foreign table update. This function can call ExplainPropertyText and related functions to add fields to the EXPLAIN output. The flag fields in es can be used to determine what to print, and the state of the ModifyTableState node can be inspected to provide run-time statistics in the EXPLAIN ANALYZE case. The first four arguments are the same as for BeginForeignModify.

If the ExplainForeignModify pointer is set to NULL, no additional information is printed during EXPLAIN.

void
ExplainDirectModify (ForeignScanState *node,
                      ExplainState *es);

Print additional EXPLAIN output for a direct modification on the remote server. This function can call ExplainPropertyText and related functions to add fields to the EXPLAIN output. The flag fields in es can be used to determine what to print, and the state of the ForeignScanState node can be inspected to provide run-time statistics in the EXPLAIN ANALYZE case.

If the ExplainDirectModify pointer is set to NULL, no additional information is printed during EXPLAIN.
55.2.7. FDW Routines for ANALYZE

```c
bool AnalyzeForeignTable (Relation relation,
    AcquireSampleRowsFunc *func,
    BlockNumber *totalpages);
```

This function is called when ANALYZE is executed on a foreign table. If the FDW can collect statistics for this foreign table, it should return `true`, and provide a pointer to a function that will collect sample rows from the table in `func`, plus the estimated size of the table in pages in `totalpages`. Otherwise, return `false`.

If the FDW does not support collecting statistics for any tables, the `AnalyzeForeignTable` pointer can be set to `NULL`.

If provided, the sample collection function must have the signature

```c
int
    AcquireSampleRowsFunc (Relation relation, int elevel,
        HeapTuple *rows, int targrows,
        double *totalrows,
        double *totaldeadrows);
```

A random sample of up to `targrows` rows should be collected from the table and stored into the caller-provided `rows` array. The actual number of rows collected must be returned. In addition, store estimates of the total numbers of live and dead rows in the table into the output parameters `totalrows` and `totaldeadrows`. (Set `totaldeadrows` to zero if the FDW does not have any concept of dead rows.)

55.2.8. FDW Routines For IMPORT FOREIGN SCHEMA

```c
List *
    ImportForeignSchema (ImportForeignSchemaStmt *stmt, Oid serverOid);
```

Obtain a list of foreign table creation commands. This function is called when executing IMPORT FOREIGN SCHEMA, and is passed the parse tree for that statement, as well as the OID of the foreign server to use. It should return a list of C strings, each of which must contain a CREATE FOREIGN TABLE command. These strings will be parsed and executed by the core server.

Within the `ImportForeignSchemaStmt` struct, `remote_schema` is the name of the remote schema from which tables are to be imported. `list_type` identifies how to filter table names: `FDW_IMPORT_SCHEMA_ALL` means that all tables in the remote schema should be imported (in this case `table_list` is empty), `FDW_IMPORT_SCHEMA_LIMIT_TO` means to include only tables listed in `table_list`, and `FDW_IMPORT_SCHEMA_EXCEPT` means to exclude the tables listed in `table_list`. `options` is a list of options used for the import process. The meanings of the options are up to the FDW. For example, an FDW could use an option to define whether the `NOT NULL` attributes of columns should be imported. These options need not have anything to do with those supported by the FDW as database object options.

The FDW may ignore the `local_schema` field of the `ImportForeignSchemaStmt`, because the core server will automatically insert that name into the parsed CREATE FOREIGN TABLE commands.

The FDW does not have to concern itself with implementing the filtering specified by `list_type` and `table_list`, either, as the core server will automatically skip any returned commands for tables.
excluded according to those options. However, it’s often useful to avoid the work of creating commands for excluded tables in the first place. The function `IsImportableForeignTable()` may be useful to test whether a given foreign-table name will pass the filter.

If the FDW does not support importing table definitions, the `ImportForeignSchema` pointer can be set to `NULL`.

### 55.2.9. FDW Routines for Parallel Execution

A `ForeignScan` node can, optionally, support parallel execution. A parallel `ForeignScan` will be executed in multiple processes and should return each row only once across all cooperating processes. To do this, processes can coordinate through fixed size chunks of dynamic shared memory. This shared memory is not guaranteed to be mapped at the same address in every process, so pointers may not be used. The following callbacks are all optional in general, but required if parallel execution is to be supported.

```c
bool IsForeignScanParallelSafe(PlannerInfo *root, RelOptInfo *rel, RangeTblEntry *rte);
```

Test whether a scan can be performed within a parallel worker. This function will only be called when the planner believes that a parallel plan might be possible, and should return true if it is safe for that scan to run within a parallel worker. This will generally not be the case if the remote data source has transaction semantics, unless the worker’s connection to the data can somehow be made to share the same transaction context as the leader.

If this callback is not defined, it is assumed that the scan must take place within the parallel leader. Note that returning true does not mean that the scan itself can be done in parallel, only that the scan can be performed within a parallel worker. Therefore, it can be useful to define this method even when parallel execution is not supported.

```c
Size EstimateDSMForeignScan(ForeignScanState *node, ParallelContext *pcxt);
```

Estimate the amount of dynamic shared memory that will be required for parallel operation. This may be higher than the amount that will actually be used, but it must not be lower. The return value is in bytes.

```c
void InitializeDSMForeignScan(ForeignScanState *node, ParallelContext *pcxt, void *coordinate);
```

Initialize the dynamic shared memory that will be required for parallel operation. `coordinate` points to an amount of allocated space equal to the return value of `EstimateDSMForeignScan`.

```c
void InitializeWorkerForeignScan(ForeignScanState *node, shm_toc *toc, void *coordinate);
```

Initialize a parallel worker’s custom state based on the shared state set up in the leader by `InitializeDSMForeignScan`. This callback is optional, and needs only be supplied if this custom path supports parallel execution.
55.3. Foreign Data Wrapper Helper Functions

Several helper functions are exported from the core server so that authors of foreign data wrappers can get easy access to attributes of FDW-related objects, such as FDW options. To use any of these functions, you need to include the header file `foreign/foreign.h` in your source file. That header also defines the struct types that are returned by these functions.

```c
ForeignDataWrapper * GetForeignDataWrapper(Oid fdwid);
```

This function returns a `ForeignDataWrapper` object for the foreign-data wrapper with the given OID. A `ForeignDataWrapper` object contains properties of the FDW (see `foreign/foreign.h` for details).

```c
ForeignServer * GetForeignServer(Oid serverid);
```

This function returns a `ForeignServer` object for the foreign server with the given OID. A `ForeignServer` object contains properties of the server (see `foreign/foreign.h` for details).

```c
UserMapping * GetUserMapping(Oid userid, Oid serverid);
```

This function returns a `UserMapping` object for the user mapping of the given role on the given server. (If there is no mapping for the specific user, it will return the mapping for PUBLIC, or throw error if there is none.) A `UserMapping` object contains properties of the user mapping (see `foreign/foreign.h` for details).

```c
ForeignTable * GetForeignTable(Oid relid);
```

This function returns a `ForeignTable` object for the foreign table with the given OID. A `ForeignTable` object contains properties of the foreign table (see `foreign/foreign.h` for details).

```c
List * GetForeignColumnOptions(Oid relid, AttrNumber attnum);
```

This function returns the per-column FDW options for the column with the given foreign table OID and attribute number, in the form of a list of `DefElem`. NIL is returned if the column has no options.

Some object types have name-based lookup functions in addition to the OID-based ones:

```c
ForeignDataWrapper * GetForeignDataWrapperByName(const char *name, bool missing_ok);
```

This function returns a `ForeignDataWrapper` object for the foreign-data wrapper with the given name. If the wrapper is not found, return NULL if missing_ok is true, otherwise raise an error.

```c
ForeignServer * GetForeignServerByName(const char *name, bool missing_ok);
```

This function returns a `ForeignServer` object for the foreign server with the given name. If the server is not found, return NULL if missing_ok is true, otherwise raise an error.
55.4. Foreign Data Wrapper Query Planning

The FDW callback functions `GetForeignRelSize`, `GetForeignPaths`, `GetForeignPlan`, `PlanForeignModify`, `GetForeignJoinPaths`, `GetForeignUpperPaths`, and `PlanDirectModify` must fit into the workings of the PostgreSQL planner. Here are some notes about what they must do.

The information in `root` and `baserel` can be used to reduce the amount of information that has to be fetched from the foreign table (and therefore reduce the cost). `baserel->baserestrictinfo` is particularly interesting, as it contains restriction quals (WHERE clauses) that should be used to filter the rows to be fetched. (The FDW itself is not required to enforce these quals, as the core executor can check them instead.) `baserel->reltarget->exprs` can be used to determine which columns need to be fetched; but note that it only lists columns that have to be emitted by the `ForeignScan` plan node, not columns that are used in qual evaluation but not output by the query.

Various private fields are available for the FDW planning functions to keep information in. Generally, whatever you store in FDW private fields should be palloc’d, so that it will be reclaimed at the end of planning.

`baserel->fdw_private` is a void pointer that is available for FDW planning functions to store information relevant to the particular foreign table. The core planner does not touch it except to initialize it to NULL when the `RelOptInfo` node is created. It is useful for passing information forward from `GetForeignRelSize` to `GetForeignPaths` and/or `GetForeignPaths` to `GetForeignPlan`, thereby avoiding recalculation.

`GetForeignPaths` can identify the meaning of different access paths by storing private information in the `fdw_private` field of `ForeignPath` nodes. `fdw_private` is declared as a List pointer, but could actually contain anything since the core planner does not touch it. However, best practice is to use a representation that’s dumpable by `nodeToString`, for use with debugging support available in the backend.

In `GetForeignPlan`, generally the passed-in target list can be copied into the plan node as-is. The passed `scan_clauses` list contains the same clauses as `baserel->baserestrictinfo`, but may be re-ordered for better execution efficiency. In simple cases the FDW can just strip `RestrictInfo` nodes from the `scan_clauses` list (using `extract_actual_clauses`) and put all the clauses into the plan node’s qual list, which means that all the clauses will be checked by the executor at run time. More complex FDWs may be able to check some of the clauses internally, in which case those clauses can be removed from the plan node’s qual list so that the executor doesn’t waste time rechecking them.

As an example, the FDW might identify some restriction clauses of the form `foreign_variable = sub_expression`, which it determines can be executed on the remote server given the locally-evaluated value of the `sub_expression`. The actual identification of such a clause should happen during `GetForeignPaths`, since it would affect the cost estimate for the path. The path’s `fdw_private` field would probably include a pointer to the identified clause’s `RestrictInfo` node. Then `GetForeignPlan` would remove that clause from `scan_clauses`, but add the `sub_expression` to `fdw_exprs` to ensure that it gets massaged into executable form. It would probably also put control information into the plan node’s `fdw_private` field to tell the execution functions what to do at run time. The query transmitted to the remote server would involve
something like WHERE foreign_variable = $1, with the parameter value obtained at run time from evaluation of the fdw_exprs expression tree.

Any clauses removed from the plan node’s qual list must instead be added to fdw_recheck_quals or rechecked by RecheckForeignScan in order to ensure correct behavior at the READ COMMITTED isolation level. When a concurrent update occurs for some other table involved in the query, the executor may need to verify that all of the original quals are still satisfied for the tuple, possibly against a different set of parameter values. Using fdw_recheck_quals is typically easier than implementing checks inside RecheckForeignScan, but this method will be insufficient when outer joins have been pushed down, since the join tuples in that case might have some fields go to NULL without rejecting the tuple entirely.

Another ForeignScan field that can be filled by FDWs is fdw_scan_tlist, which describes the tuples returned by the FDW for this plan node. For simple foreign table scans this can be set to NIL, implying that the returned tuples have the row type declared for the foreign table. A non-NIL value must be a target list (list of TargetEntries) containing Vars and/or expressions representing the returned columns. This might be used, for example, to show that the FDW has omitted some columns that it noticed won’t be needed for the query. Also, if the FDW can compute expressions used by the query more cheaply than can be done locally, it could add those expressions to fdw_scan_tlist.

The FDW should always construct at least one path that depends only on the table’s restriction clauses. In join queries, it might also choose to construct path(s) that depend on join clauses, for example foreign_variable = local_variable. Such clauses will not be found in baserel->baserestrictinfo but must be sought in the relation’s join lists. A path using such a clause is called a “parameterized path”. It must identify the other relations used in the selected join clause(s) with a suitable value of param_info; use get_baserel_parampathinfo to compute that value. In GetForeignPlan, the local_variable portion of the join clause would be added to fdw_exprs, and then at run time the case works the same as for an ordinary restriction clause.

If an FDW supports remote joins, GetForeignJoinPaths should produce ForeignPaths for potential remote joins in much the same way as GetForeignPaths works for base tables. Information about the intended join can be passed forward to GetForeignPlan in the same ways described above. However, baserestrictinfo is not relevant for join relations; instead, the relevant join clauses for a particular join are passed to GetForeignJoinPaths as a separate parameter (extra->restrictlist).

An FDW might additionally support direct execution of some plan actions that are above the level of scans and joins, such as grouping or aggregation. To offer such options, the FDW should generate paths and insert them into the appropriate upper relation. For example, a path representing remote aggregation should be inserted into the UPPERREL_GROUP_AGG relation, using add_path. This path will be compared on a cost basis with local aggregation performed by reading a simple scan path for the foreign relation (note that such a path must also be supplied, else there will be an error at plan time). If the remote-aggregation path wins, which it usually would, it will be converted into a plan in the usual way, by calling GetForeignPlan. The recommended place to generate such paths is in the GetForeignUpperPaths callback function, which is called for each upper relation (i.e., each post-scan/join processing step), if all the base relations of the query come from the same FDW.

PlanForeignModify and the other callbacks described in Section 55.2.4 are designed around the assumption that the foreign relation will be scanned in the usual way and then individual row updates will be driven by a local ModifyTable plan node. This approach is necessary for the general case where an update requires reading local tables as well as foreign tables. However, if the operation could be executed entirely by the foreign server, the FDW could generate a path representing that and insert it into the UPPERREL_FINAL upper relation, where it would compete against the ModifyTable
approach. This approach could also be used to implement remote `SELECT FOR UPDATE`, rather than using the row locking callbacks described in Section 55.2.5. Keep in mind that a path inserted into `UPPERREL_FINAL` is responsible for implementing all behavior of the query.

When planning an `UPDATE` or `DELETE`, `PlanForeignModify` and `PlanDirectModify` can look up the `RelOptInfo` struct for the foreign table and make use of the `baserel->fdw_private` data previously created by the scan-planning functions. However, in `INSERT` the target table is not scanned so there is no `RelOptInfo` for it. The list returned by `PlanForeignModify` has the same restrictions as the `fdw_private` list of a `ForeignScan` plan node, that is it must contain only structures that `copyObject` knows how to copy.

`INSERT` with an `ON CONFLICT` clause does not support specifying the conflict target, as unique constraints or exclusion constraints on remote tables are not locally known. This in turn implies that `ON CONFLICT DO UPDATE` is not supported, since the specification is mandatory there.

55.5. Row Locking in Foreign Data Wrappers

If an FDW’s underlying storage mechanism has a concept of locking individual rows to prevent concurrent updates of those rows, it is usually worthwhile for the FDW to perform row-level locking with as close an approximation as practical to the semantics used in ordinary PostgreSQL tables. There are multiple considerations involved in this.

One key decision to be made is whether to perform early locking or late locking. In early locking, a row is locked when it is first retrieved from the underlying store, while in late locking, the row is locked only when it is known that it needs to be locked. (The difference arises because some rows may be discarded by locally-checked restriction or join conditions.) Early locking is much simpler and avoids extra round trips to a remote store, but it can cause locking of rows that need not have been locked, resulting in reduced concurrency or even unexpected deadlocks. Also, late locking is only possible if the row to be locked can be uniquely re-identified later. Preferably the row identifier should identify a specific version of the row, as PostgreSQL TIDs do.

By default, PostgreSQL ignores locking considerations when interfacing to FDWs, but an FDW can perform early locking without any explicit support from the core code. The API functions described in Section 55.2.5, which were added in PostgreSQL 9.5, allow an FDW to use late locking if it wishes.

An additional consideration is that in `READ COMMITTED` isolation mode, PostgreSQL may need to re-check restriction and join conditions against an updated version of some target tuple. Rechecking join conditions requires re-obtaining copies of the non-target rows that were previously joined to the target tuple. When working with standard PostgreSQL tables, this is done by including the TIDs of the non-target tables in the column list projected through the join, and then re-fetching non-target rows when required. This approach keeps the join data set compact, but it requires inexpensive re-fetch capability, as well as a TID that can uniquely identify the row version to be re-fetched. By default, therefore, the approach used with foreign tables is to include a copy of the entire row fetched from a foreign table in the column list projected through the join. This puts no special demands on the FDW but can result in reduced performance of merge and hash joins. An FDW that is capable of meeting the re-fetch requirements can choose to do it the first way.

For an `UPDATE` or `DELETE` on a foreign table, it is recommended that the `ForeignScan` operation on the target table perform early locking on the rows that it fetches, perhaps via the equivalent of `SELECT FOR UPDATE`. An FDW can detect whether a table is an `UPDATE/DELETE` target at plan time by comparing its `relid` to `root->parse->resultRelation`, or at execution time by using `ExecRelationIsTargetRelation()`. An alternative possibility is to perform late locking within the `ExecForeignUpdate` or `ExecForeignDelete` callback, but no special support is provided for this.
For foreign tables that are specified to be locked by a `SELECT FOR UPDATE/SHARE` command, the `ForeignScan` operation can again perform early locking by fetching tuples with the equivalent of `SELECT FOR UPDATE/SHARE`. To perform late locking instead, provide the callback functions defined in Section 55.2.5. In `GetForeignRowMarkType`, select rowmark option `ROW_MARK_EXCLUSIVE`, `ROW_MARK_NOKEYEXCLUSIVE`, `ROW_MARK_SHARE`, or `ROW_MARK_KEYSHARE` depending on the requested lock strength. (The core code will act the same regardless of which of these four options you choose.) Elsewhere, you can detect whether a foreign table was specified to be locked by this type of command by using `get_plan_rowmark` at plan time, or `ExecFindRowMark` at execution time; you must check not only whether a non-null rowmark struct is returned, but that its `strength` field is not `LCS_NONE`.

Lastly, for foreign tables that are used in an `UPDATE`, `DELETE` or `SELECT FOR UPDATE/SHARE` command but are not specified to be row-locked, you can override the default choice to copy entire rows by having `GetForeignRowMarkType` select option `ROW_MARK_REFERENCE` when it sees lock strength `LCS_NONE`. This will cause `RefetchForeignRow` to be called with that value for `markType`; it should then re-fetch the row without acquiring any new lock. (If you have a `GetForeignRowMarkType` function but don’t wish to re-fetch unlocked rows, select option `ROW_MARK_COPY` for `LCS_NONE`.)

See `src/include/nodes/lockoptions.h`, the comments for `RowMarkType` and `PlanRowMark` in `src/include/nodes/plannodes.h`, and the comments for `ExecRowMark` in `src/include/nodes/execnodes.h` for additional information.
Chapter 56. Writing A Table Sampling Method

PostgreSQL’s implementation of the TABLESAMPLE clause supports custom table sampling methods, in addition to the BERNOULLI and SYSTEM methods that are required by the SQL standard. The sampling method determines which rows of the table will be selected when the TABLESAMPLE clause is used.

At the SQL level, a table sampling method is represented by a single SQL function, typically implemented in C, having the signature

\[\text{method_name} (\text{internal}) \quad \text{RETURNS} \quad \text{tsm_handler}\]

The name of the function is the same method name appearing in the TABLESAMPLE clause. The \text{internal} argument is a dummy (always having value zero) that simply serves to prevent this function from being called directly from a SQL command. The result of the function must be a palloc’d struct of type \text{TsmRoutine}, which contains pointers to support functions for the sampling method. These support functions are plain C functions and are not visible or callable at the SQL level. The support functions are described in Section 56.1.

In addition to function pointers, the \text{TsmRoutine} struct must provide these additional fields:

\text{List} \quad \text{parameterTypes}

This is an OID list containing the data type OIDs of the parameter(s) that will be accepted by the TABLESAMPLE clause when this sampling method is used. For example, for the built-in methods, this list contains a single item with value \text{FLOAT4OID}, which represents the sampling percentage. Custom sampling methods can have more or different parameters.

\text{bool} \quad \text{repeatable_across_queries}

If true, the sampling method can deliver identical samples across successive queries, if the same parameters and \text{REPEATABLE} seed value are supplied each time and the table contents have not changed. When this is false, the \text{REPEATABLE} clause is not accepted for use with the sampling method.

\text{bool} \quad \text{repeatable_across_scans}

If true, the sampling method can deliver identical samples across successive scans in the same query (assuming unchanging parameters, seed value, and snapshot). When this is false, the planner will not select plans that would require scanning the sampled table more than once, since that might result in inconsistent query output.

The \text{TsmRoutine} struct type is declared in src/include/access/tsmapi.h, which see for additional details.

The table sampling methods included in the standard distribution are good references when trying to write your own. Look into the src/backend/access/tablesample subdirectory of the source tree for the built-in sampling methods, and into the contrib subdirectory for add-on methods.

56.1. Sampling Method Support Functions

The TSM handler function returns a palloc’d \text{TsmRoutine} struct containing pointers to the support functions described below. Most of the functions are required, but some are optional, and those point-
Chapter 56. Writing A Table Sampling Method

void
SampleScanGetSampleSize (PlannerInfo *root,
    RelOptInfo *baserel,
    List *paramexprs,
    BlockNumber *pages,
    double *tuples);

This function is called during planning. It must estimate the number of relation pages that will be read during a sample scan, and the number of tuples that will be selected by the scan. (For example, these might be determined by estimating the sampling fraction, and then multiplying the baserel->pages and baserel->tuples numbers by that, being sure to round the results to integral values.) The paramexprs list holds the expression(s) that are parameters to the TABLESAMPLE clause. It is recommended to use estimate_expression_value() to try to reduce these expressions to constants, if their values are needed for estimation purposes; but the function must provide size estimates even if they cannot be reduced, and it should not fail even if the values appear invalid (remember that they’re only estimates of what the run-time values will be). The pages and tuples parameters are outputs.

void
InitSampleScan (SampleScanState *node,
        int eflags);

Initialize for execution of a SampleScan plan node. This is called during executor startup. It should perform any initialization needed before processing can start. The SampleScanState node has already been created, but its tsm_state field is NULL. The InitSampleScan function can palloc whatever internal state data is needed by the sampling method, and store a pointer to it in node->tsm_state. Information about the table to scan is accessible through other fields of the SampleScanState node (but note that the node->ss.ss_currentScanDesc scan descriptor is not set up yet). eflags contains flag bits describing the executor’s operating mode for this plan node.

When (eflags & EXEC_FLAG_EXPLAIN_ONLY) is true, the scan will not actually be performed, so this function should only do the minimum required to make the node state valid for EXPLAIN and EndSampleScan.

This function can be omitted (set the pointer to NULL), in which case BeginSampleScan must perform all initialization needed by the sampling method.

void
BeginSampleScan (SampleScanState *node,
        Datum *params,
        int nparams,
        uint32 seed);

Begin execution of a sampling scan. This is called just before the first attempt to fetch a tuple, and may be called again if the scan needs to be restarted. Information about the table to scan is accessible through fields of the SampleScanState node (but note that the node->ss.ss_currentScanDesc scan descriptor is not set up yet). The params array, of length nparams, contains the values of the parameters supplied in the TABLESAMPLE clause. These will have the number and types specified in the sampling method’s parameterTypes list, and have been checked to not be null. seed contains a seed to use for any random numbers generated within the sampling method; it is either a hash derived from the REPEATABLE value if one was given, or the result of random() if not.
This function may adjust the fields node->use_bulkread and node->use_pagemode. If node->use_bulkread is true, which it is by default, the scan will use a buffer access strategy that encourages recycling buffers after use. It might be reasonable to set this to false if the scan will visit only a small fraction of the table’s pages. If node->use_pagemode is true, which it is by default, the scan will perform visibility checking in a single pass for all tuples on each visited page. It might be reasonable to set this to false if the scan will select only a small fraction of the tuples on each visited page. That will result in fewer tuple visibility checks being performed, though each one will be more expensive because it will require more locking.

If the sampling method is marked repeatable_across_scans, it must be able to select the same set of tuples during a rescan as it did originally, that is a fresh call of BeginSampleScan must lead to selecting the same tuples as before (if the TABLESAMPLE parameters and seed don’t change).

```
BlockNumber
NextSampleBlock (SampleScanState *node);
```

Returns the block number of the next page to be scanned, or InvalidBlockNumber if no pages remain to be scanned.

This function can be omitted (set the pointer to NULL), in which case the core code will perform a sequential scan of the entire relation. Such a scan can use synchronized scanning, so that the sampling method cannot assume that the relation pages are visited in the same order on each scan.

```
OffsetNumber
NextSampleTuple (SampleScanState *node, 
                 BlockNumber blockno, 
                 OffsetNumber maxoffset);
```

Returns the offset number of the next tuple to be sampled on the specified page, or InvalidOffsetNumber if no tuples remain to be sampled. maxoffset is the largest offset number in use on the page.

**Note:** NextSampleTuple is not explicitly told which of the offset numbers in the range 1 .. maxoffset actually contain valid tuples. This is not normally a problem since the core code ignores requests to sample missing or invisible tuples; that should not result in any bias in the sample. However, if necessary, the function can examine node->ss.ss_currentScanDesc->rs_vistuples[] to identify which tuples are valid and visible. (This requires node->use_pagemode to be true.)

**Note:** NextSampleTuple must not assume that blockno is the same page number returned by the most recent NextSampleBlock call. It was returned by some previous NextSampleBlock call, but the core code is allowed to call NextSampleBlock in advance of actually scanning pages, so as to support prefetching. It is OK to assume that once sampling of a given page begins, successive NextSampleTuple calls all refer to the same page until InvalidOffsetNumber is returned.

```
void
EndSampleScan (SampleScanState *node);
```

End the scan and release resources. It is normally not important to release palloc’d memory, but any externally-visible resources should be cleaned up. This function can be omitted (set the pointer to NULL) in the common case where no such resources exist.
Chapter 57. Writing A Custom Scan Provider

PostgreSQL supports a set of experimental facilities which are intended to allow extension modules to add new scan types to the system. Unlike a foreign data wrapper, which is only responsible for knowing how to scan its own foreign tables, a custom scan provider can provide an alternative method of scanning any relation in the system. Typically, the motivation for writing a custom scan provider will be to allow the use of some optimization not supported by the core system, such as caching or some form of hardware acceleration. This chapter outlines how to write a new custom scan provider.

Implementing a new type of custom scan is a three-step process. First, during planning, it is necessary to generate access paths representing a scan using the proposed strategy. Second, if one of those access paths is selected by the planner as the optimal strategy for scanning a particular relation, the access path must be converted to a plan. Finally, it must be possible to execute the plan and generate the same results that would have been generated for any other access path targeting the same relation.

57.1. Creating Custom Scan Paths

A custom scan provider will typically add paths for a base relation by setting the following hook, which is called after the core code has generated all the access paths it can for the relation (except for Gather paths, which are made after this call so that they can use partial paths added by the hook):

```
typedef void (*set_rel_pathlist_hook_type) (PlannerInfo *root, RelOptInfo *rel, Index rti, RangeTblEntry *rte);
extern PGDLLIMPORT set_rel_pathlist_hook_type set_rel_pathlist_hook;
```

Although this hook function can be used to examine, modify, or remove paths generated by the core system, a custom scan provider will typically confine itself to generating CustomPath objects and adding them to rel using add_path. The custom scan provider is responsible for initializing the CustomPath object, which is declared like this:

```
typedef struct CustomPath
{
    Path path;
    uint32 flags;
    List *custom_paths;
    List *custom_private;
    const CustomPathMethods *methods;
} CustomPath;
```

path must be initialized as for any other path, including the row-count estimate, start and total cost, and sort ordering provided by this path. flags is a bit mask, which should include CUSTOMPATH_SUPPORT_BACKWARD_SCAN if the custom path can support a backward scan and CUSTOMPATH_SUPPORT_MARK_RESTORE if it can support mark and restore. Both capabilities are optional. An optional custom_paths is a list of Path nodes used by this custom-path node; these will be transformed into Plan nodes by planner. custom_private can be used to store the custom
Chapter 57. Writing A Custom Scan Provider

path’s private data. Private data should be stored in a form that can be handled by nodeToString, so that debugging routines that attempt to print the custom path will work as designed. methods must point to a (usually statically allocated) object implementing the required custom path methods, of which there is currently only one.

A custom scan provider can also provide join paths. Just as for base relations, such a path must produce the same output as would normally be produced by the join it replaces. To do this, the join provider should set the following hook, and then within the hook function, create CustomPath path(s) for the join relation.

typedef void (*set_join_pathlist_hook_type) (PlannerInfo *root, RelOptInfo *joinrel, RelOptInfo *outerrel, RelOptInfo *innerrel, JoinType jointype, JoinPathExtraData *extra);

extern PGDLLIMPORT set_join_pathlist_hook_type set_join_pathlist_hook;

This hook will be invoked repeatedly for the same join relation, with different combinations of inner and outer relations; it is the responsibility of the hook to minimize duplicated work.

57.1.1. Custom Scan Path Callbacks

Plan *(*PlanCustomPath) (PlannerInfo *root, RelOptInfo *rel, CustomPath *best_path, List *tlist, List *clauses, List *custom_plans);

Convert a custom path to a finished plan. The return value will generally be a CustomScan object, which the callback must allocate and initialize. See Section 57.2 for more details.

57.2. Creating Custom Scan Plans

A custom scan is represented in a finished plan tree using the following structure:

typedef struct CustomScan
{
    Scan         scan;
    uint32       flags;
    List         *custom_plans;
    List         *custom_exprs;
    List         *custom_private;
    List         *custom_scan_tlist;
    Bitmapset    *custom_relids;
    const CustomScanMethods  *methods;
} CustomScan;

scan must be initialized as for any other scan, including estimated costs, target lists, qualifications, and so on. flags is a bit mask with the same meaning as in CustomPath. custom_plans can be
used to store child Plan nodes. custom_exprs should be used to store expression trees that will need to be fixed up by setrefs.c and subselect.c, while custom_private should be used to store other private data that is only used by the custom scan provider itself. custom_scan_tlist can be NIL when scanning a base relation, indicating that the custom scan returns scan tuples that match the base relation’s row type. Otherwise it is a target list describing the actual scan tuples. custom_scan_tlist must be provided for joins, and could be provided for scans if the custom scan provider can compute some non-Var expressions. custom_relids is set by the core code to the set of relations (range table indexes) that this scan node handles; except when this scan is replacing a join, it will have only one member. methods must point to a (usually statically allocated) object implementing the required custom scan methods, which are further detailed below.

When a CustomScan scans a single relation, scan.scanrelid must be the range table index of the table to be scanned. When it replaces a join, scan.scanrelid should be zero.

Plan trees must be able to be duplicated using copyObject, so all the data stored within the “custom” fields must consist of nodes that that function can handle. Furthermore, custom scan providers cannot substitute a larger structure that embeds a CustomScan for the structure itself, as would be possible for a CustomPath or CustomScanState.

57.2.1. Custom Scan Plan Callbacks

Node *(CreateCustomScanState) (CustomScan *cscan);

Allocate a CustomScanState for this CustomScan. The actual allocation will often be larger than required for an ordinary CustomScanState, because many providers will wish to embed that as the first field of a larger structure. The value returned must have the node tag and methods set appropriately, but other fields should be left as zeroes at this stage; after ExecInitCustomScan performs basic initialization, the BeginCustomScan callback will be invoked to give the custom scan provider a chance to do whatever else is needed.

57.3. Executing Custom Scans

When a CustomScan is executed, its execution state is represented by a CustomScanState, which is declared as follows:

typedef struct CustomScanState
{
    ScanState ss;
    uint32 flags;
    const CustomExecMethods *methods;
} CustomScanState;

ss is initialized as for any other scan state, except that if the scan is for a join rather than a base relation, ss.ss_currentRelation is left NULL. flags is a bit mask with the same meaning as in CustomPath and CustomScan. methods must point to a (usually statically allocated) object implementing the required custom scan state methods, which are further detailed below. Typically, a CustomScanState, which need not support copyObject, will actually be a larger structure embedding the above as its first member.
57.3.1. Custom Scan Execution Callbacks

void (\*BeginCustomScan) (CustomScanState *node,
    EState *estate,
    int eflags);

Complete initialization of the supplied CustomScanState. Standard fields have been initialized by ExecInitCustomScan, but any private fields should be initialized here.

TupleTableSlot *(\*ExecCustomScan) (CustomScanState *node);

Fetch the next scan tuple. If any tuples remain, it should fill ps_ResultTupleSlot with the next tuple in the current scan direction, and then return the tuple slot. If not, NULL or an empty slot should be returned.

void (\*EndCustomScan) (CustomScanState *node);

Clean up any private data associated with the CustomScanState. This method is required, but it does not need to do anything if there is no associated data or it will be cleaned up automatically.

void (\*ReScanCustomScan) (CustomScanState *node);

Rewind the current scan to the beginning and prepare to rescan the relation.

void (\*MarkPosCustomScan) (CustomScanState *node);

Save the current scan position so that it can subsequently be restored by the RestrPosCustomScan callback. This callback is optional, and need only be supplied if the CUSTOMPATH_SUPPORT_MARK_RESTORE flag is set.

void (\*RestrPosCustomScan) (CustomScanState *node);

Restore the previous scan position as saved by the MarkPosCustomScan callback. This callback is optional, and need only be supplied if the CUSTOMPATH_SUPPORT_MARK_RESTORE flag is set.

Size (\*EstimateDSMCustomScan) (CustomScanState *node,
    ParallelContext *pcxt);

Estimate the amount of dynamic shared memory that will be required for parallel operation. This may be higher than the amount that will actually be used, but it must not be lower. The return value is in bytes. This callback is optional, and need only be supplied if this custom scan provider supports parallel execution.

void (\*InitializeDSMCustomScan) (CustomScanState *node,
    ParallelContext *pcxt,
    void *coordinate);

Initialize the dynamic shared memory that will be required for parallel operation; coordinate points to an amount of allocated space equal to the return value of EstimateDSMCustomScan. This callback is optional, and need only be supplied if this custom scan provider supports parallel execution.

void (\*InitializeWorkerCustomScan) (CustomScanState *node,
    shm_toc *toc,
Chapter 57. Writing A Custom Scan Provider

void *coordinate);

Initialize a parallel worker’s custom state based on the shared state set up in the leader by InitializeDSMCustomScan. This callback is optional, and needs only be supplied if this custom path supports parallel execution.

void (*ExplainCustomScan) (CustomScanState *node,
                             List *ancestors,
                             ExplainState *es);

Output additional information for EXPLAIN of a custom-scan plan node. This callback is optional. Common data stored in the ScanState, such as the target list and scan relation, will be shown even without this callback, but the callback allows the display of additional, private state.
Chapter 58. Genetic Query Optimizer

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58.1. Query Handling as a Complex Optimization Problem

Among all relational operators the most difficult one to process and optimize is the join. The number of possible query plans grows exponentially with the number of joins in the query. Further optimization effort is caused by the support of a variety of join methods (e.g., nested loop, hash join, merge join in PostgreSQL) to process individual joins and a diversity of indexes (e.g., B-tree, hash, GiST and GIN in PostgreSQL) as access paths for relations.

The normal PostgreSQL query optimizer performs a near-exhaustive search over the space of alternative strategies. This algorithm, first introduced in IBM’s System R database, produces a near-optimal join order, but can take an enormous amount of time and memory space when the number of joins in the query grows large. This makes the ordinary PostgreSQL query optimizer inappropriate for queries that join a large number of tables.

The Institute of Automatic Control at the University of Mining and Technology, in Freiberg, Germany, encountered some problems when it wanted to use PostgreSQL as the backend for a decision support knowledge based system for the maintenance of an electrical power grid. The DBMS needed to handle large join queries for the inference machine of the knowledge based system. The number of joins in these queries made using the normal query optimizer infeasible.

In the following we describe the implementation of a genetic algorithm to solve the join ordering problem in a manner that is efficient for queries involving large numbers of joins.

58.2. Genetic Algorithms

The genetic algorithm (GA) is a heuristic optimization method which operates through randomized search. The set of possible solutions for the optimization problem is considered as a population of individuals. The degree of adaptation of an individual to its environment is specified by its fitness.

The coordinates of an individual in the search space are represented by chromosomes, in essence a set of character strings. A gene is a subsection of a chromosome which encodes the value of a single parameter being optimized. Typical encodings for a gene could be binary or integer.

Through simulation of the evolutionary operations recombination, mutation, and selection new generations of search points are found that show a higher average fitness than their ancestors.
According to the comp.ai.genetic FAQ it cannot be stressed too strongly that a GA is not a pure random search for a solution to a problem. A GA uses stochastic processes, but the result is distinctly non-random (better than random).

Figure 58-1. Structured Diagram of a Genetic Algorithm

```
+=========================================+
| P(t) generation of ancestors at a time t |
| P'(t) generation of descendants at a time t |
+=========================================+
```

58.3. Genetic Query Optimization (GEQO) in PostgreSQL

The GEQO module approaches the query optimization problem as though it were the well-known traveling salesman problem (TSP). Possible query plans are encoded as integer strings. Each string represents the join order from one relation of the query to the next. For example, the join tree

```
  /
 /  \
/    2
/  3
```

is encoded by the integer string '4-1-3-2', which means, first join relation '4' and '1', then '3', and then '2', where 1, 2, 3, 4 are relation IDs within the PostgreSQL optimizer.

Specific characteristics of the GEQO implementation in PostgreSQL are:

- Usage of a steady state GA (replacement of the least fit individuals in a population, not whole-generational replacement) allows fast convergence towards improved query plans. This is essential for query handling with reasonable time;
Chapter 58. Genetic Query Optimizer

- Usage of *edge recombination crossover* which is especially suited to keep edge losses low for the solution of the TSP by means of a GA;
- Mutation as genetic operator is deprecated so that no repair mechanisms are needed to generate legal TSP tours.

Parts of the GEQO module are adapted from D. Whitley’s Genitor algorithm.

The GEQO module allows the PostgreSQL query optimizer to support large join queries effectively through non-exhaustive search.

### 58.3.1. Generating Possible Plans with GEQO

The GEQO planning process uses the standard planner code to generate plans for scans of individual relations. Then join plans are developed using the genetic approach. As shown above, each candidate join plan is represented by a sequence in which to join the base relations. In the initial stage, the GEQO code simply generates some possible join sequences at random. For each join sequence considered, the standard planner code is invoked to estimate the cost of performing the query using that join sequence. (For each step of the join sequence, all three possible join strategies are considered; and all the initially-determined relation scan plans are available. The estimated cost is the cheapest of these possibilities.) Join sequences with lower estimated cost are considered “more fit” than those with higher cost. The genetic algorithm discards the least fit candidates. Then new candidates are generated by combining genes of more-fit candidates — that is, by using randomly-chosen portions of known low-cost join sequences to create new sequences for consideration. This process is repeated until a preset number of join sequences have been considered; then the best one found at any time during the search is used to generate the finished plan.

This process is inherently nondeterministic, because of the randomized choices made during both the initial population selection and subsequent “mutation” of the best candidates. To avoid surprising changes of the selected plan, each run of the GEQO algorithm restarts its random number generator with the current `geqo_seed` parameter setting. As long as `geqo_seed` and the other GEQO parameters are kept fixed, the same plan will be generated for a given query (and other planner inputs such as statistics). To experiment with different search paths, try changing `geqo_seed`.

### 58.3.2. Future Implementation Tasks for PostgreSQL GEQO

Work is still needed to improve the genetic algorithm parameter settings. In file `src/backend/optimizer/geqo/geqo_main.c`, routines `gimme_pool_size` and `gimme_number_generations`, we have to find a compromise for the parameter settings to satisfy two competing demands:

- Optimality of the query plan
- Computing time

In the current implementation, the fitness of each candidate join sequence is estimated by running the standard planner’s join selection and cost estimation code from scratch. To the extent that different candidates use similar sub-sequences of joins, a great deal of work will be repeated. This could be made significantly faster by retaining cost estimates for sub-joins. The problem is to avoid expending unreasonable amounts of memory on retaining that state.
At a more basic level, it is not clear that solving query optimization with a GA algorithm designed for TSP is appropriate. In the TSP case, the cost associated with any substring (partial tour) is independent of the rest of the tour, but this is certainly not true for query optimization. Thus it is questionable whether edge recombination crossover is the most effective mutation procedure.

58.4. Further Reading

The following resources contain additional information about genetic algorithms:

- Evolutionary Computation and its application to art and design², by Craig Reynolds
- Fundamentals of Database Systems
- *The design and implementation of the POSTGRES query optimizer*

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¹. http://www.aip.de/~ast/EvolCompFAQ/
Chapter 59. Index Access Method Interface Definition

This chapter defines the interface between the core PostgreSQL system and index access methods, which manage individual index types. The core system knows nothing about indexes beyond what is specified here, so it is possible to develop entirely new index types by writing add-on code.

All indexes in PostgreSQL are what are known technically as secondary indexes; that is, the index is physically separate from the table file that it describes. Each index is stored as its own physical relation and so is described by an entry in the pg_class catalog. The contents of an index are entirely under the control of its index access method. In practice, all index access methods divide indexes into standard-size pages so that they can use the regular storage manager and buffer manager to access the index contents. (All the existing index access methods furthermore use the standard page layout described in Section 65.6, and most use the same format for index tuple headers; but these decisions are not forced on an access method.)

An index is effectively a mapping from some data key values to tuple identifiers, or TIDs, of row versions (tuples) in the index’s parent table. A TID consists of a block number and an item number within that block (see Section 65.6). This is sufficient information to fetch a particular row version from the table. Indexes are not directly aware that under MVCC, there might be multiple extant versions of the same logical row; to an index, each tuple is an independent object that needs its own index entry. Thus, an update of a row always creates all-new index entries for the row, even if the key values did not change. (HOT tuples are an exception to this statement; but indexes do not deal with those, either.) Index entries for dead tuples are reclaimed (by vacuuming) when the dead tuples themselves are reclaimed.

59.1. Basic API Structure for Indexes

Each index access method is described by a row in the pg_am system catalog. The pg_am entry specifies a name and a handler function for the access method. These entries can be created and deleted using the CREATE ACCESS METHOD and DROP ACCESS METHOD SQL commands.

An index access method handler function must be declared to accept a single argument of type internal and to return the pseudo-type index_am_handler. The argument is a dummy value that simply serves to prevent handler functions from being called directly from SQL commands. The result of the function must be a malloc’d struct of type IndexAmRoutine, which contains everything that the core code needs to know to make use of the index access method. The IndexAmRoutine struct, also called the access method’s API struct, includes fields specifying assorted fixed properties of the access method, such as whether it can support multicolumn indexes. More importantly, it contains pointers to support functions for the access method, which do all of the real work to access indexes. These support functions are plain C functions and are not visible or callable at the SQL level. The support functions are described in Section 59.2.

The structure IndexAmRoutine is defined thus:

typedef struct IndexAmRoutine
{
    NodeTag    type;
}
Chapter 59. Index Access Method Interface Definition

/*
 * Total number of strategies (operators) by which we can traverse/search
 * this AM. Zero if AM does not have a fixed set of strategy assignments.
 */
uint16 amstrategies;
/* total number of support functions that this AM uses */
uint16 amsupport;
/* does AM support ORDER BY indexed column’s value? */
bool amcanorder;
/* does AM support ORDER BY result of an operator on indexed column? */
bool amcanorderbyop;
/* does AM support backward scanning? */
bool amcanbackward;
/* does AM support UNIQUE indexes? */
bool amcanunique;
/* does AM support multi-column indexes? */
bool amcanmulticol;
/* does AM require scans to have a constraint on the first index column? */
bool amoptionalkey;
/* does AM handle ScalarArrayOpExpr quals? */
bool amsearcharray;
/* does AM handle IS NULL/IS NOT NULL quals? */
bool amsearchnulls;
/* can index storage data type differ from column data type? */
bool amstorage;
/* can an index of this type be clustered on? */
bool amclusterable;
/* does AM handle predicate locks? */
bool ampredlocks;
/* type of data stored in index, or InvalidOid if variable */
Oid amkeytype;

/* interface functions */
ambuild_function ambuild;
ambuildempty_function ambuildempty;
amininsert_function amininsert;
ambulkdelete_function ambulkdelete;
amvacuumcleanup_function amvacuumcleanup;
amcanreturn_function amcanreturn;
/* can be NULL */
amcostestimate_function amcostestimate;
amoptions_function amoptions;
amproperty_function amproperty;
/* can be NULL */
amvalidate_function amvalidate;
ambeginscan_function ambeginscan;
amrescan_function amrescan;
amgettuple_function amgettuple;
/* can be NULL */
amgetbitmap_function amgetbitmap;
/* can be NULL */
amendscan_function amendscan;
ammarkpos_function ammarkpos;
/* can be NULL */
amrestrpos_function amrestrpos;
/* can be NULL */
} IndexAmRoutine;

To be useful, an index access method must also have one or more operator families and operator classes defined in pg_opfamily, pg_opclass, pg_amop, and pg_amproc. These entries allow the planner to determine what kinds of query qualifications can be used with indexes of this access.
method. Operator families and classes are described in Section 36.14, which is prerequisite material for reading this chapter.

An individual index is defined by a `pg_class` entry that describes it as a physical relation, plus a `pg_index` entry that shows the logical content of the index — that is, the set of index columns it has and the semantics of those columns, as captured by the associated operator classes. The index columns (key values) can be either simple columns of the underlying table or expressions over the table rows. The index access method normally has no interest in where the index key values come from (it is always handed precomputed key values) but it will be very interested in the operator class information in `pg_index`. Both of these catalog entries can be accessed as part of the `Relation` data structure that is passed to all operations on the index.

Some of the flag fields of `IndexAmRoutine` have nonobvious implications. The requirements of `amcanunique` are discussed in Section 59.5. The `amcanmulticol` flag asserts that the access method supports multicol indexes, while `amoptionalkey` asserts that it allows scans where no indexable restriction clause is given for the first index column. When `amcanmulticol` is false, `amoptionalkey` essentially says whether the access method supports full-index scans without any restriction clause. Access methods that support multiple index columns must support scans that omit restrictions on any or all of the columns after the first; however they are permitted to require some restriction to appear for the first index column, and this is signaled by setting `amoptionalkey` false. One reason that an index AM might set `amoptionalkey` false is if it doesn’t index null values. Since most indexable operators are strict and hence cannot return true for null inputs, it is at first sight attractive to not store index entries for null values: they could never be returned by an index scan anyway. However, this argument fails when an index scan has no restriction clause for a given index column. In practice this means that indexes that have `amoptionalkey` true must index nulls, since the planner might decide to use such an index with no scan keys at all. A related restriction is that an index access method that supports multiple index columns must support indexing null values in columns after the first, because the planner will assume the index can be used for queries that do not restrict these columns. For example, consider an index on (a,b) and a query with `WHERE a = 4`. The system will assume the index can be used to scan for rows with `a = 4`, which is wrong if the index omits rows where `b` is null. It is, however, OK to omit rows where the first indexed column is null. An index access method that does index nulls may also set `amsearchnulls`, indicating that it supports `IS NULL` and `IS NOT NULL` clauses as search conditions.

### 59.2. Index Access Method Functions

The index construction and maintenance functions that an index access method must provide in `IndexAmRoutine` are:

```c
IndexBuildResult * ambuild (Relation heapRelation,
                             Relation indexRelation,
                             IndexInfo *indexInfo);
```

Build a new index. The index relation has been physically created, but is empty. It must be filled in with whatever fixed data the access method requires, plus entries for all tuples already existing in the table. Ordinarily the `ambuild` function will call `IndexBuildHeapScan()` to scan the table for existing tuples and compute the keys that need to be inserted into the index. The function must return a `palloc`d struct containing statistics about the new index.

```c
void ambuildempty (Relation indexRelation);
```
Build an empty index, and write it to the initialization fork (INIT_FORKNUM) of the given relation. This method is called only for unlogged indexes; the empty index written to the initialization fork will be copied over the main relation fork on each server restart.

```c
bool aminsert (Relation indexRelation,
    Datum *values,
    bool *isnull,
    ItemPointer heap_tid,
    Relation heapRelation,
    IndexUniqueCheck checkUnique);
```

Insert a new tuple into an existing index. The values and isnull arrays give the key values to be indexed, and heap_tid is the TID to be indexed. If the access method supports unique indexes (its amcanunique flag is true) then checkUnique indicates the type of uniqueness check to perform. This varies depending on whether the unique constraint is deferrable; see Section 59.5 for details. Normally the access method only needs the heapRelation parameter when performing uniqueness checking (since then it will have to look into the heap to verify tuple liveness).

The function’s Boolean result value is significant only when checkUnique is UNIQUE_CHECK_PARTIAL. In this case a TRUE result means the new entry is known unique, whereas FALSE means it might be non-unique (and a deferred uniqueness check must be scheduled). For other cases a constant FALSE result is recommended.

Some indexes might not index all tuples. If the tuple is not to be indexed, aminsert should just return without doing anything.

```c
IndexBulkDeleteResult * ambulkdelete (IndexVacuumInfo *info,
    IndexBulkDeleteResult *stats,
    IndexBulkDeleteCallback callback,
    void *callback_state);
```

Delete tuple(s) from the index. This is a “bulk delete” operation that is intended to be implemented by scanning the whole index and checking each entry to see if it should be deleted. The passed-in callback function must be called, in the style callback(TID, callback_state) returns bool, to determine whether any particular index entry, as identified by its referenced TID, is to be deleted. Must return either NULL or a palloc’d struct containing statistics about the effects of the deletion operation. It is OK to return NULL if no information needs to be passed on to amvacuumcleanup.

Because of limited maintenance_work_mem, ambulkdelete might need to be called more than once when many tuples are to be deleted. The stats argument is the result of the previous call for this index (it is NULL for the first call within a VACUUM operation). This allows the AM to accumulate statistics across the whole operation. Typically, ambulkdelete will modify and return the same struct if the passed stats is not null.

```c
IndexBulkDeleteResult * amvacuumcleanup (IndexVacuumInfo *info,
    IndexBulkDeleteResult *stats);
```

Clean up after a VACUUM operation (zero or more ambulkdelete calls). This does not have to do anything beyond returning index statistics, but it might perform bulk cleanup such as reclaiming empty index pages. stats is whatever the last ambulkdelete call returned, or NULL if ambulkdelete was not called because no tuples needed to be deleted. If the result is not NULL it must be a palloc’d
struct. The statistics it contains will be used to update \texttt{pg\_class}, and will be reported by \texttt{VACUUM} if \texttt{VERBOSE} is given. It is OK to return \texttt{NULL} if the index was not changed at all during the \texttt{VACUUM} operation, but otherwise correct stats should be returned.

As of PostgreSQL 8.4, \texttt{amvacuumcleanup} will also be called at completion of an \texttt{ANALYZE} operation. In this case \texttt{stats} is always \texttt{NULL} and any return value will be ignored. This case can be distinguished by checking \texttt{info->analyze\_only}. It is recommended that the access method do nothing except post-insert cleanup in such a call, and that only in an autovacuum worker process.

```c
bool amcanreturn (Relation indexRelation, int attno);
```

Check whether the index can support index-only scans on the given column, by returning the indexed column values for an index entry in the form of an \texttt{IndexTuple}. The attribute number is 1-based, i.e. the first column’s attno is 1. Returns TRUE if supported, else FALSE. If the access method does not support index-only scans at all, the \texttt{amcanreturn} field in its \texttt{IndexAmRoutine} struct can be set to \texttt{NULL}.

```c
void amcostestimate (PlannerInfo *root,
    IndexPath *path,
    double loop_count,
    Cost *indexStartupCost,
    Cost *indexTotalCost,
    Selectivity *indexSelectivity,
    double *indexCorrelation);
```

Estimate the costs of an index scan. This function is described fully in Section 59.6, below.

```c
bytea * amoptions (ArrayType *reloptions,
    bool validate);
```

Parse and validate the reloptions array for an index. This is called only when a non-null reloptions array exists for the index. \texttt{reloptions} is a text array containing entries of the form \texttt{name=value}. The function should construct a \texttt{bytea} value, which will be copied into the \texttt{rd\_options} field of the index’s relcache entry. The data contents of the \texttt{bytea} value are open for the access method to define; most of the standard access methods use struct \texttt{StdRdOptions}. When \texttt{validate} is true, the function should report a suitable error message if any of the options are unrecognized or have invalid values; when \texttt{validate} is false, invalid entries should be silently ignored. (\texttt{validate} is false when loading options already stored in \texttt{pg\_catalog}; an invalid entry could only be found if the access method has changed its rules for options, and in that case ignoring obsolete entries is appropriate.) It is OK to return \texttt{NULL} if default behavior is wanted.

```c
bool amproperty (Oid index\_oid, int attno,
    IndexAMProperty prop, const char *propname,
    bool *res, bool *isnull);
```

The \texttt{amproperty} method allows index access methods to override the default behavior of \texttt{pg\_index\_column\_has\_property} and related functions. If the access method does not have any special behavior for index property inquiries, the \texttt{amproperty} field in its \texttt{IndexAmRoutine} struct can be set to \texttt{NULL}. Otherwise, the \texttt{amproperty} method will be called with \texttt{index\_oid} and \texttt{attno} both zero for \texttt{pg\_indexam\_has\_property} calls, or with \texttt{index\_oid} valid and \texttt{attno} zero.
for pg_index_has_property calls, or with index_oid valid and attno greater than zero for pg_index_column_has_property calls. prop is an enum value identifying the property being tested, while proppname is the original property name string. If the core code does not recognize the property name then prop is AMPROP_UNKNOWN. Access methods can define custom property names by checking proppname for a match (use pg_strcasecmp to match, for consistency with the core code); for names known to the core code, it’s better to inspect prop. If the amproperty method returns true then it has determined the property test result: it must set *res to the boolean value to return, or set *isnull to true to return a NULL. (Both of the referenced variables are initialized to false before the call.) If the amproperty method returns false then the core code will proceed with its normal logic for determining the property test result.

Access methods that support ordering operators should implement AMPROP_DISTANCE_ORDERABLE property testing, as the core code does not know how to do that and will return NULL. It may also be advantageous to implement AMPROP_RETURNABLE testing, if that can be done more cheaply than by opening the index and calling amcanreturn, which is the core code’s default behavior. The default behavior should be satisfactory for all other standard properties.

```c
bool
amvalidate (Oid opclassoid);
```

Validate the catalog entries for the specified operator class, so far as the access method can reasonably do that. For example, this might include testing that all required support functions are provided. The amvalidate function must return false if the opclass is invalid. Problems should be reported with ereport messages.

The purpose of an index, of course, is to support scans for tuples matching an indexable WHERE condition, often called a qualifier or scan key. The semantics of index scanning are described more fully in Section 59.3, below. An index access method can support “plain” index scans, “bitmap” index scans, or both. The scan-related functions that an index access method must or may provide are:

```c
IndexScanDesc
ambeginScan (Relation indexRelation,
              int nkeys,
              int norderbys);
```

Prepare for an index scan. The nkeys and norderbys parameters indicate the number of quals and ordering operators that will be used in the scan; these may be useful for space allocation purposes. Note that the actual values of the scan keys aren’t provided yet. The result must be a pal-lc’d struct. For implementation reasons the index access method must create this struct by calling RelationGetIndexScan(). In most cases ambeginScan does little beyond making that call and perhaps acquiring locks; the interesting parts of index-scan startup are in amrescan.

```c
void
amrescan (IndexScanDesc scan,
           ScanKey keys,
           int nkeys,
           ScanKey orderbys,
           int norderbys);
```

Start or restart an index scan, possibly with new scan keys. (To restart using previously-passed keys, NULL is passed for keys and/or orderbys.) Note that it is not allowed for the number of keys or order-by operators to be larger than what was passed to ambeginScan. In practice the restart feature is used when a new outer tuple is selected by a nested-loop join and so a new key comparison value is needed, but the scan key structure remains the same.
boolean
amgettuple (IndexScanDesc scan,
   ScanDirection direction);

Fetch the next tuple in the given scan, moving in the given direction (forward or backward in the index). Returns TRUE if a tuple was obtained, FALSE if no matching tuples remain. In the TRUE case the tuple TID is stored into the scan structure. Note that “success” means only that the index contains an entry that matches the scan keys, not that the tuple necessarily still exists in the heap or will pass the caller’s snapshot test. On success, amgettuple must also set scan->xs_recheck to TRUE or FALSE. FALSE means it is certain that the index entry matches the scan keys. TRUE means this is not certain, and the conditions represented by the scan keys must be rechecked against the heap tuple after fetching it. This provision supports “lossy” index operators. Note that rechecking will extend only to the scan conditions; a partial index predicate (if any) is never rechecked by amgettuple callers.

If the index supports index-only scans (i.e., amcanreturn returns TRUE for it), then on success the AM must also check scan->xs_want_itup, and if that is true it must return the original indexed data for the index entry, in the form of an IndexTuple pointer stored at scan->xs_itup, with tuple descriptor scan->xs_itupdesc. (Management of the data referenced by the pointer is the access method’s responsibility. The data must remain good at least until the next amgettuple, amrescan, or amendscan call for the scan.)

The amgettuple function need only be provided if the access method supports “plain” index scans. If it doesn’t, the amgettuple field in its IndexAmRoutine struct must be set to NULL.

int64
amgetbitmap (IndexScanDesc scan,
   TIDBitmap *tbm);

Fetch all tuples in the given scan and add them to the caller-supplied TIDBitmap (that is, OR the set of tuple IDs into whatever set is already in the bitmap). The number of tuples fetched is returned (this might be just an approximate count, for instance some AMs do not detect duplicates). While inserting tuple IDs into the bitmap, amgetbitmap can indicate that rechecking of the scan conditions is required for specific tuple IDs. This is analogous to the xs_recheck output parameter of amgettuple. Note: in the current implementation, support for this feature is conflated with support for lossy storage of the bitmap itself, and therefore callers recheck both the scan conditions and the partial index predicate (if any) for recheckable tuples. That might not always be true, however. amgetbitmap and amgettuple cannot be used in the same index scan; there are other restrictions too when using amgetbitmap, as explained in Section 59.3.

The amgetbitmap function need only be provided if the access method supports “bitmap” index scans. If it doesn’t, the amgetbitmap field in its IndexAmRoutine struct must be set to NULL.

void
amendscan (IndexScanDesc scan);

End a scan and release resources. The scan struct itself should not be freed, but any locks or pins taken internally by the access method must be released, as well as any other memory allocated by ambeginscan and other scan-related functions.

void
ammarkpos (IndexScanDesc scan);

Mark current scan position. The access method need only support one remembered scan position per scan.
The `ammarkpos` function need only be provided if the access method supports ordered scans. If it doesn’t, the `ammarkpos` field in its `IndexAmRoutine` struct may be set to NULL.

```c
void
amrestrpos (IndexScanDesc scan);
```

Restore the scan to the most recently marked position.

The `amrestrpos` function need only be provided if the access method supports ordered scans. If it doesn’t, the `amrestrpos` field in its `IndexAmRoutine` struct may be set to NULL.

### 59.3. Index Scanning

In an index scan, the index access method is responsible for regurgitating the TIDs of all the tuples it has been told about that match the `scan keys`. The access method is not involved in actually fetching those tuples from the index’s parent table, nor in determining whether they pass the scan’s time qualification test or other conditions.

A scan key is the internal representation of a `WHERE` clause of the form `index_key operator constant`, where the index key is one of the columns of the index and the operator is one of the members of the operator family associated with that index column. An index scan has zero or more scan keys, which are implicitly ANDed — the returned tuples are expected to satisfy all the indicated conditions.

The access method can report that the index is lossy, or requires rechecks, for a particular query. This implies that the index scan will return all the entries that pass the scan key, plus possibly additional entries that do not. The core system’s index-scan machinery will then apply the index conditions again to the heap tuple to verify whether or not it really should be selected. If the recheck option is not specified, the index scan must return exactly the set of matching entries.

Note that it is entirely up to the access method to ensure that it correctly finds all and only the entries passing all the given scan keys. Also, the core system will simply hand off all the `WHERE` clauses that match the index keys and operator families, without any semantic analysis to determine whether they are redundant or contradictory. As an example, given `WHERE x > 4 AND x > 14` where `x` is a b-tree indexed column, it is left to the b-tree `amrescan` function to realize that the first scan key is redundant and can be discarded. The extent of preprocessing needed during `amrescan` will depend on the extent to which the index access method needs to reduce the scan keys to a “normalized” form.

Some access methods return index entries in a well-defined order, others do not. There are actually two different ways that an access method can support sorted output:

- Access methods that always return entries in the natural ordering of their data (such as btree) should set `amcanorder` to true. Currently, such access methods must use btree-compatible strategy numbers for their equality and ordering operators.
- Access methods that support ordering operators should set `amcanorderbyop` to true. This indicates that the index is capable of returning entries in an order satisfying `ORDER BY index_key operator constant`. Scan modifiers of that form can be passed to `amrescan` as described previously.

The `amgettuple` function has a `direction` argument, which can be either `ForwardScanDirection` (the normal case) or `BackwardScanDirection`. If the first call after `amrescan` specifies `BackwardScanDirection`, then the set of matching index entries is to be
scanned back-to-front rather than in the normal front-to-back direction, so amgettuple must return the last matching tuple in the index, rather than the first one as it normally would. (This will only occur for access methods that set amcanorder to true.) After the first call, amgettuple must be prepared to advance the scan in either direction from the most recently returned entry. (But if amcanbackward is false, all subsequent calls will have the same direction as the first one.)

Access methods that support ordered scans must support “marking” a position in a scan and later returning to the marked position. The same position might be restored multiple times. However, only one position need be remembered per scan; a new ammarkpos call overrides the previously marked position. An access method that does not support ordered scans need not provide ammarkpos and amrestrpos functions in IndexAmRoutine; set those pointers to NULL instead.

Both the scan position and the mark position (if any) must be maintained consistently in the face of concurrent insertions or deletions in the index. It is OK if a freshly-inserted entry is not returned by a scan that would have found the entry if it had existed when the scan started, or for the scan to return such an entry upon rescanning or backing up even though it had not been returned the first time through. Similarly, a concurrent delete might or might not be reflected in the results of a scan. What is important is that insertions or deletions not cause the scan to miss or multiply return entries that were not themselves being inserted or deleted.

If the index stores the original indexed data values (and not some lossy representation of them), it is useful to support index-only scans, in which the index returns the actual data not just the TID of the heap tuple. This will only avoid I/O if the visibility map shows that the TID is on an all-visible page; else the heap tuple must be visited anyway to check MVCC visibility. But that is no concern of the access method’s.

Instead of using amgettuple, an index scan can be done with amgetbitmap to fetch all tuples in one call. This can be noticeably more efficient than amgettuple because it allows avoiding lock/unlock cycles within the access method. In principle amgetbitmap should have the same effects as repeated amgettuple calls, but we impose several restrictions to simplify matters. First of all, amgetbitmap returns all tuples at once and marking or restoring scan positions isn’t supported. Secondly, the tuples are returned in a bitmap which doesn’t have any specific ordering, which is why amgetbitmap doesn’t take a direction argument. (Ordering operators will never be supplied for such a scan, either.) Also, there is no provision for index-only scans with amgetbitmap, since there is no way to return the contents of index tuples. Finally, amgetbitmap does not guarantee any locking of the returned tuples, with implications spelled out in Section 59.4.

Note that it is permitted for an access method to implement only amgetbitmap and not amgettuple, or vice versa, if its internal implementation is unsuited to one API or the other.

### 59.4. Index Locking Considerations

Index access methods must handle concurrent updates of the index by multiple processes. The core PostgreSQL system obtains AccessShareLock on the index during an index scan, and RowExclusiveLock when updating the index (including plain VACUUM). Since these lock types do not conflict, the access method is responsible for handling any fine-grained locking it might need. An exclusive lock on the index as a whole will be taken only during index creation, destruction, or REINDEX.

Building an index type that supports concurrent updates usually requires extensive and subtle analysis of the required behavior. For the b-tree and hash index types, you can read about the design decisions involved in src/backend/access/nbtree/README and src/backend/access/hash/README.
Aside from the index’s own internal consistency requirements, concurrent updates create issues about consistency between the parent table (the heap) and the index. Because PostgreSQL separates accesses and updates of the heap from those of the index, there are windows in which the index might be inconsistent with the heap. We handle this problem with the following rules:

- A new heap entry is made before making its index entries. (Therefore a concurrent index scan is likely to fail to see the heap entry. This is okay because the index reader would be uninterested in an uncommitted row anyway. But see Section 59.5.)
- When a heap entry is to be deleted (by VACUUM), all its index entries must be removed first.
- An index scan must maintain a pin on the index page holding the item last returned by amgettuple, and ambulkdelete cannot delete entries from pages that are pinned by other backends. The need for this rule is explained below.

Without the third rule, it is possible for an index reader to see an index entry just before it is removed by VACUUM, and then to arrive at the corresponding heap entry after that was removed by VACUUM. This creates no serious problems if that item number is still unused when the reader reaches it, since an empty item slot will be ignored by heap_fetch(). But what if a third backend has already re-used the item slot for something else? When using an MVCC-compliant snapshot, there is no problem because the new occupant of the slot is certain to be too new to pass the snapshot test. However, with a non-MVCC-compliant snapshot (such as SnapshotAny), it would be possible to accept and return a row that does not in fact match the scan keys. We could defend against this scenario by requiring the scan keys to be rechecked against the heap row in all cases, but that is too expensive. Instead, we use a pin on an index page as a proxy to indicate that the reader might still be “in flight” from the index entry to the matching heap entry. Making ambulkdelete block on such a pin ensures that VACUUM cannot delete the heap entry before the reader is done with it. This solution costs little in run time, and adds blocking overhead only in the rare cases where there actually is a conflict.

This solution requires that index scans be “synchronous”: we have to fetch each heap tuple immediately after scanning the corresponding index entry. This is expensive for a number of reasons. An “asynchronous” scan in which we collect many TIDs from the index, and only visit the heap tuples sometime later, requires much less index locking overhead and can allow a more efficient heap access pattern. Per the above analysis, we must use the synchronous approach for non-MVCC-compliant snapshots, but an asynchronous scan is workable for a query using an MVCC snapshot.

In an amgetbitmap index scan, the access method does not keep an index pin on any of the returned tuples. Therefore it is only safe to use such scans with MVCC-compliant snapshots.

When the ampredlocks flag is not set, any scan using that index access method within a serializable transaction will acquire a nonblocking predicate lock on the full index. This will generate a read-write conflict with the insert of any tuple into that index by a concurrent serializable transaction. If certain patterns of read-write conflicts are detected among a set of concurrent serializable transactions, one of those transactions may be canceled to protect data integrity. When the flag is set, it indicates that the index access method implements finer-grained predicate locking, which will tend to reduce the frequency of such transaction cancellations.

### 59.5. Index Uniqueness Checks

PostgreSQL enforces SQL uniqueness constraints using unique indexes, which are indexes that disallow multiple entries with identical keys. An access method that supports this feature sets amcanunique true. (At present, only b-tree supports it.)
Because of MVCC, it is always necessary to allow duplicate entries to exist physically in an index: the entries might refer to successive versions of a single logical row. The behavior we actually want to enforce is that no MVCC snapshot could include two rows with equal index keys. This breaks down into the following cases that must be checked when inserting a new row into a unique index:

- If a conflicting valid row has been deleted by the current transaction, it’s okay. (In particular, since an UPDATE always deletes the old row version before inserting the new version, this will allow an UPDATE on a row without changing the key.)
- If a conflicting row has been inserted by an as-yet-uncommitted transaction, the would-be inserter must wait to see if that transaction commits. If it rolls back then there is no conflict. If it commits without deleting the conflicting row again, there is a uniqueness violation. (In practice we just wait for the other transaction to end and then redo the visibility check in toto.)
- Similarly, if a conflicting valid row has been deleted by an as-yet-uncommitted transaction, the would-be inserter must wait for that transaction to commit or abort, and then repeat the test.

Furthermore, immediately before reporting a uniqueness violation according to the above rules, the access method must recheck the liveness of the row being inserted. If it is committed dead then no violation should be reported. (This case cannot occur during the ordinary scenario of inserting a row that’s just been created by the current transaction. It can happen during \texttt{CREATE UNIQUE INDEX CONCURRENTLY}, however.)

We require the index access method to apply these tests itself, which means that it must reach into the heap to check the commit status of any row that is shown to have a duplicate key according to the index contents. This is without a doubt ugly and non-modular, but it saves redundant work: if we did a separate probe then the index lookup for a conflicting row would be essentially repeated while finding the place to insert the new row’s index entry. What’s more, there is no obvious way to avoid race conditions unless the conflict check is an integral part of insertion of the new index entry.

If the unique constraint is deferrable, there is additional complexity: we need to be able to insert an index entry for a new row, but defer any uniqueness-violation error until end of statement or even later. To avoid unnecessary repeat searches of the index, the index access method should do a preliminary uniqueness check during the initial insertion. If this shows that there is definitely no conflicting live tuple, we are done. Otherwise, we schedule a recheck to occur when it is time to enforce the constraint. If, at the time of the recheck, both the inserted tuple and some other tuple with the same key are live, then the error must be reported. (Note that for this purpose, “live” actually means “any tuple in the index entry’s HOT chain is live”. To implement this, the \texttt{aminsert} function is passed a \texttt{checkUnique} parameter having one of the following values:

- \texttt{UNIQUE_CHECK_NO} indicates that no uniqueness checking should be done (this is not a unique index).
- \texttt{UNIQUE_CHECK_YES} indicates that this is a non-deferrable unique index, and the uniqueness check must be done immediately, as described above.
- \texttt{UNIQUE_CHECK_PARTIAL} indicates that the unique constraint is deferrable. PostgreSQL will use this mode to insert each row’s index entry. The access method must allow duplicate entries into the index, and report any potential duplicates by returning \texttt{FALSE} from \texttt{aminsert}. For each row for which \texttt{FALSE} is returned, a deferred recheck will be scheduled.

The access method must identify any rows which might violate the unique constraint, but it is not an error for it to report false positives. This allows the check to be done without waiting for other transactions to finish; conflicts reported here are not treated as errors and will be rechecked later, by which time they may no longer be conflicts.
Chapter 59. Index Access Method Interface Definition

- **UNIQUE_CHECK_EXISTING** indicates that this is a deferred recheck of a row that was reported as a potential uniqueness violation. Although this is implemented by calling `aminsert`, the access method must not insert a new index entry in this case. The index entry is already present. Rather, the access method must check to see if there is another live index entry. If so, and if the target row is also still live, report error.

It is recommended that in a **UNIQUE_CHECK_EXISTING** call, the access method further verify that the target row actually does have an existing entry in the index, and report error if not. This is a good idea because the index tuple values passed to `aminsert` will have been recomputed. If the index definition involves functions that are not really immutable, we might be checking the wrong area of the index. Checking that the target row is found in the recheck verifies that we are scanning for the same tuple values as were used in the original insertion.

### 59.6. Index Cost Estimation Functions

The `amcostestimate` function is given information describing a possible index scan, including lists of WHERE and ORDER BY clauses that have been determined to be usable with the index. It must return estimates of the cost of accessing the index and the selectivity of the WHERE clauses (that is, the fraction of parent-table rows that will be retrieved during the index scan). For simple cases, nearly all the work of the cost estimator can be done by calling standard routines in the optimizer; the point of having an `amcostestimate` function is to allow index access methods to provide index-type-specific knowledge, in case it is possible to improve on the standard estimates.

Each `amcostestimate` function must have the signature:

```c
void amcostestimate (PlannerInfo *root,
                      IndexPath *path,
                      double loop_count,
                      Cost *indexStartupCost,
                      Cost *indexTotalCost,
                      Selectivity *indexSelectivity,
                      double *indexCorrelation);
```

The first three parameters are inputs:

- **root**
  
The planner’s information about the query being processed.

- **path**
  
The index access path being considered. All fields except cost and selectivity values are valid.

- **loop_count**
  
The number of repetitions of the index scan that should be factored into the cost estimates. This will typically be greater than one when considering a parameterized scan for use in the inside of a nestloop join. Note that the cost estimates should still be for just one scan; a larger `loop_count` means that it may be appropriate to allow for some caching effects across multiple scans.

The last four parameters are pass-by-reference outputs:

- **indexStartupCost**
  
- **indexTotalCost**
  
- **indexSelectivity**
  
- **indexCorrelation**
Chapter 59. Index Access Method Interface Definition

*indexStartupCost

Set to cost of index start-up processing

*indexTotalCost

Set to total cost of index processing

*indexSelectivity

Set to index selectivity

*indexCorrelation

Set to correlation coefficient between index scan order and underlying table’s order

Note that cost estimate functions must be written in C, not in SQL or any available procedural language, because they must access internal data structures of the planner/optimizer.

The index access costs should be computed using the parameters used by src/backend/optimizer/path/costsize.c: a sequential disk block fetch has cost seq_page_cost, a nonsequential fetch has cost random_page_cost, and the cost of processing one index row should usually be taken as cpu_index_tuple_cost. In addition, an appropriate multiple of cpu_operator_cost should be charged for any comparison operators invoked during index processing (especially evaluation of the indexquals themselves).

The access costs should include all disk and CPU costs associated with scanning the index itself, but not the costs of retrieving or processing the parent-table rows that are identified by the index.

The “start-up cost” is the part of the total scan cost that must be expended before we can begin to fetch the first row. For most indexes this can be taken as zero, but an index type with a high start-up cost might want to set it nonzero.

The indexSelectivity should be set to the estimated fraction of the parent table rows that will be retrieved during the index scan. In the case of a lossy query, this will typically be higher than the fraction of rows that actually pass the given qual conditions.

The indexCorrelation should be set to the correlation (ranging between -1.0 and 1.0) between the index order and the table order. This is used to adjust the estimate for the cost of fetching rows from the parent table.

When loop_count is greater than one, the returned numbers should be averages expected for any one scan of the index.

Cost Estimation

A typical cost estimator will proceed as follows:

1. Estimate and return the fraction of parent-table rows that will be visited based on the given qual conditions. In the absence of any index-type-specific knowledge, use the standard optimizer function clauselist_selectivity():

   ```c
   *indexSelectivity = clauselist_selectivity(root, path->indexquals,
      path->indexinfo->rel->relid,
      JOIN.Inner, NULL);
   ```

2. Estimate the number of index rows that will be visited during the scan. For many index types this is the same as indexSelectivity times the number of rows in the index, but it might be more. (Note that the index’s size in pages and rows is available from the path->indexinfo struct.)

3. Estimate the number of index pages that will be retrieved during the scan. This might be just indexSelectivity times the index’s size in pages.
4. Compute the index access cost. A generic estimator might do this:

```c
/*
 * Our generic assumption is that the index pages will be read
 * sequentially, so they cost seq_page_cost each, not random_page_cost.
 * Also, we charge for evaluation of the indexquals at each index row.
 * All the costs are assumed to be paid incrementally during the scan.
 */

cost_qual_eval(&index_qual_cost, path->indexquals, root);
*indexStartupCost = index_qual_cost.startup;
*indexTotalCost = seq_page_cost * numIndexPages +
    (cpu_index_tuple_cost + index_qual_cost.per_tuple) * numIndexTuples;

However, the above does not account for amortization of index reads across repeated index scans.

5. Estimate the index correlation. For a simple ordered index on a single field, this can be retrieved
from pg_statistic. If the correlation is not known, the conservative estimate is zero (no correla-
tion).

Examples of cost estimator functions can be found in `src/backend/utils/adt/selfuncs.c`. 
Chapter 60. Generic WAL Records

Although all built-in WAL-logged modules have their own types of WAL records, there is also a generic WAL record type, which describes changes to pages in a generic way. This is useful for extensions that provide custom access methods, because they cannot register their own WAL redo routines.

The API for constructing generic WAL records is defined in `access/generic_xlog.h` and implemented in `access/transam/generic_xlog.c`.

To perform a WAL-logged data update using the generic WAL record facility, follow these steps:

1. `state = GenericXLogStart(relation)` — start construction of a generic WAL record for the given relation.
2. `page = GenericXLogRegisterBuffer(state, buffer, flags)` — register a buffer to be modified within the current generic WAL record. This function returns a pointer to a temporary copy of the buffer's page, where modifications should be made. (Do not modify the buffer's contents directly.) The third argument is a bitmask of flags applicable to the operation. Currently the only such flag is `GENERIC_XLOG_FULL_IMAGE`, which indicates that a full-page image rather than a delta update should be included in the WAL record. Typically this flag would be set if the page is new or has been rewritten completely. `GenericXLogRegisterBuffer` can be repeated if the WAL-logged action needs to modify multiple pages.
3. Apply modifications to the page images obtained in the previous step.
4. `GenericXLogFinish(state)` — apply the changes to the buffers and emit the generic WAL record.

WAL record construction can be canceled between any of the above steps by calling `GenericXLogAbort(state)`. This will discard all changes to the page image copies.

Please note the following points when using the generic WAL record facility:

- No direct modifications of buffers are allowed! All modifications must be done in copies acquired from `GenericXLogRegisterBuffer()`. In other words, code that makes generic WAL records should never call `BufferGetPage()` for itself. However, it remains the caller's responsibility to pin/umpin and lock/unlock the buffers at appropriate times. Exclusive lock must be held on each target buffer from before `GenericXLogRegisterBuffer()` until after `GenericXLogFinish()`.
- Registrations of buffers (step 2) and modifications of page images (step 3) can be mixed freely, i.e., both steps may be repeated in any sequence. Keep in mind that buffers should be registered in the same order in which locks are to be obtained on them during replay.
- The maximum number of buffers that can be registered for a generic WAL record is `MAX_GENERIC_XLOG_PAGES`. An error will be thrown if this limit is exceeded.
- Generic WAL assumes that the pages to be modified have standard layout, and in particular that there is no useful data between `pd_lower` and `pd_upper`.
- Since you are modifying copies of buffer pages, `GenericXLogStart()` does not start a critical section. Thus, you can safely do memory allocation, error throwing, etc. between `GenericXLogStart()` and `GenericXLogFinish()`. The only actual critical section is present
inside GenericXLogFinish(). There is no need to worry about calling GenericXLogAbort() during an error exit, either.

- GenericXLogFinish() takes care of marking buffers dirty and setting their LSNs. You do not need to do this explicitly.

- For unlogged relations, everything works the same except that no actual WAL record is emitted. Thus, you typically do not need to do any explicit checks for unlogged relations.

- The generic WAL redo function will acquire exclusive locks to buffers in the same order as they were registered. After redoing all changes, the locks will be released in the same order.

- If GENERIC_XLOG_FULL_IMAGE is not specified for a registered buffer, the generic WAL record contains a delta between the old and the new page images. This delta is based on byte-by-byte comparison. This is not very compact for the case of moving data within a page, and might be improved in the future.


Chapter 61. GiST Indexes

61.1. Introduction

GiST stands for Generalized Search Tree. It is a balanced, tree-structured access method, that acts as a base template in which to implement arbitrary indexing schemes. B-trees, R-trees and many other indexing schemes can be implemented in GiST.

One advantage of GiST is that it allows the development of custom data types with the appropriate access methods, by an expert in the domain of the data type, rather than a database expert.

Some of the information here is derived from the University of California at Berkeley’s GiST Indexing Project web site and Marcel Kornacker’s thesis, Access Methods for Next-Generation Database Systems. The GiST implementation in PostgreSQL is primarily maintained by Teodor Sigaev and Oleg Bartunov, and there is more information on their web site.

61.2. Built-in Operator Classes

The core PostgreSQL distribution includes the GiST operator classes shown in Table 61-1. (Some of the optional modules described in Appendix F provide additional GiST operator classes.)

Table 61-1. Built-in GiST Operator Classes

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
<th>Ordering Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>box_ops</td>
<td>box</td>
<td>&amp;&amp; &amp;&amp;&gt; &amp;&amp;&lt; &amp;&amp;&lt;</td>
<td></td>
</tr>
<tr>
<td>circle_ops</td>
<td>circle</td>
<td>&amp;&amp; &amp;&amp;&gt; &amp;&amp;&lt; &amp;&amp;&lt;</td>
<td></td>
</tr>
<tr>
<td>inet_ops</td>
<td>inet, cidr</td>
<td>&amp;&amp; &gt;&gt; &gt;&gt;= &gt; &gt; = &lt;=</td>
<td>&lt;= &gt;</td>
</tr>
<tr>
<td>point_ops</td>
<td>point</td>
<td>&gt;&gt;^ &lt;&lt; &lt;@ &lt;@</td>
<td>&lt;= &gt;</td>
</tr>
<tr>
<td>poly_ops</td>
<td>polygon</td>
<td>&amp;&amp; &amp;&amp;&gt; &amp;&amp;&lt; &amp;&amp;&lt;</td>
<td></td>
</tr>
<tr>
<td>range_ops</td>
<td>any range type</td>
<td>&amp;&amp; &amp;&amp;&gt; &amp;&amp;&lt; &amp;&amp;&lt;</td>
<td></td>
</tr>
<tr>
<td>tsquery_ops</td>
<td>tsquery</td>
<td>&lt;@ &lt;@</td>
<td>&lt;= &gt;</td>
</tr>
<tr>
<td>tsvector_ops</td>
<td>tsvector</td>
<td>&lt;@ &lt;@</td>
<td>&lt;= &gt;</td>
</tr>
</tbody>
</table>

1. http://gist.cs.berkeley.edu/
For historical reasons, the \texttt{inet\_ops} operator class is not the default class for types \texttt{inet} and \texttt{cidr}. To use it, mention the class name in \texttt{CREATE INDEX}, for example

\begin{verbatim}
CREATE INDEX ON my_table USING GIST (my_inet_column inet_ops);
\end{verbatim}

\section*{61.3. Extensibility}

Traditionally, implementing a new index access method meant a lot of difficult work. It was necessary to understand the inner workings of the database, such as the lock manager and Write-Ahead Log. The GiST interface has a high level of abstraction, requiring the access method implementer only to implement the semantics of the data type being accessed. The GiST layer itself takes care of concurrency, logging and searching the tree structure.

This extensibility should not be confused with the extensibility of the other standard search trees in terms of the data they can handle. For example, PostgreSQL supports extensible B-trees and hash indexes. That means that you can use PostgreSQL to build a B-tree or hash over any data type you want. But B-trees only support range predicates (\texttt{<}, \texttt{=}, \texttt{>}), and hash indexes only support equality queries.

So if you index, say, an image collection with a PostgreSQL B-tree, you can only issue queries such as “is imagex equal to imagey”, “is imagex less than imagey” and “is imagex greater than imagey”. Depending on how you define “equals”, “less than” and “greater than” in this context, this could be useful. However, by using a GiST based index, you could create ways to ask domain-specific questions, perhaps “find all images of horses” or “find all over-exposed images”.

All it takes to get a GiST access method up and running is to implement several user-defined methods, which define the behavior of keys in the tree. Of course these methods have to be pretty fancy to support fancy queries, but for all the standard queries (B-trees, R-trees, etc.) they’re relatively straightforward. In short, GiST combines extensibility along with generality, code reuse, and a clean interface.

There are seven methods that an index operator class for GiST must provide, and two that are optional. Correctness of the index is ensured by proper implementation of the \texttt{same}, \texttt{consistent} and \texttt{union} methods, while efficiency (size and speed) of the index will depend on the \texttt{penalty} and \texttt{picksplit} methods. The remaining two basic methods are \texttt{compress} and \texttt{decompress}, which allow an index to have internal tree data of a different type than the data it indexes. The leaves are to be of the indexed data type, while the other tree nodes can be of any C struct (but you still have to follow PostgreSQL data type rules here, see about \texttt{varlena} for variable sized data). If the tree’s internal data type exists at the SQL level, the \texttt{STORAGE} option of the \texttt{CREATE OPERATOR CLASS} command can be used. The optional eighth method is \texttt{distance}, which is needed if the operator class wishes to support ordered scans (nearest-neighbor searches). The optional ninth method \texttt{fetch} is needed if the operator class wishes to support index-only scans.

\begin{verbatim}
consistent
\end{verbatim}

Given an index entry \( p \) and a query value \( q \), this function determines whether the index entry is “consistent” with the query; that is, could the predicate “\texttt{indexed\_column indexable\_operator q}” be true for any row represented by the index entry? For a leaf index entry this is equivalent to testing the indexable condition, while for an internal tree node this determines whether it is necessary to scan the subtree of the index represented by the tree node. When the result is \texttt{true}, a \texttt{recheck} flag must also be returned. This indicates whether the predicate is certainly true or only possibly true. If \texttt{recheck = false} then the index has tested
the predicate condition exactly, whereas if `recheck = true` the row is only a candidate match. In that case the system will automatically evaluate the `indexable_operator` against the actual row value to see if it is really a match. This convention allows GiST to support both lossless and lossy index structures.

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_consistent(internal, data_type, smallint, oid, internal)
RETURNS bool
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:

```c
Datum
my_consistent(PG_FUNCTION_ARGS)
{
    GISTENTRY *entry = (GISTENTRY *) PG_GETARG_POINTER(0);
    data_type *query = PG_GETARG_DATA_TYPE_P(1);
    StrategyNumber strategy = (StrategyNumber) PG_GETARG_UINT16(2);
    /* Oid subtype = PG_GETARG_OID(3); */
    bool *recheck = (bool *) PG_GETARG_POINTER(4);
    data_type *key = DatumGetDataType(entry->key);
    bool retval;

    /*
    * determine return value as a function of strategy, key and query.
    *
    * Use GIST_LEAF(entry) to know where you're called in the index tree,
    * which comes handy when supporting the = operator for example (you could
    * check for non empty union() in non-leaf nodes and equality in leaf
    * nodes).
    */

    *recheck = true;  /* or false if check is exact */

    PG_RETURN_BOOL(retval);
}
```

Here, `key` is an element in the index and `query` the value being looked up in the index. The `StrategyNumber` parameter indicates which operator of your operator class is being applied — it matches one of the operator numbers in the `CREATE OPERATOR CLASS` command.

Depending on which operators you have included in the class, the data type of `query` could vary with the operator, since it will be whatever type is on the righthand side of the operator, which might be different from the indexed data type appearing on the lefthand side. (The above code skeleton assumes that only one type is possible; if not, fetching the `query` argument value would have to depend on the operator.) It is recommended that the SQL declaration of the `consistent` function use the opclass’s indexed data type for the `query` argument, even though the actual type might be something else depending on the operator.

```sql
CREATE OR REPLACE FUNCTION my_union(internal, internal)
RETURNS storage_type
```

This method consolidates information in the tree. Given a set of entries, this function generates a new index entry that represents all the given entries.
Chapter 61. GiST Indexes

AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:

PG_FUNCTION_INFO_V1(my_union);

Datum
my_union(PG_FUNCTION_ARGS)
{
    GistEntryVector *entryvec = (GistEntryVector *) PG_GETARG_POINTER(0);
    GISTENTRY *ent = entryvec->vector;
    data_type *out,
        *tmp,
        *old;
    int numranges,
        i = 0;

    numranges = entryvec->n;
    tmp = DatumGetDataType(ent[0].key);
    out = tmp;

    if (numranges == 1)
    {
        out = data_type_deep_copy(tmp);

        PG_RETURN_DATA_TYPE_P(out);
    }

    for (i = 1; i < numranges; i++)
    {
        old = out;
        tmp = DatumGetDataType(ent[i].key);
        out = my_union_implementation(out, tmp);
    }

    PG_RETURN_DATA_TYPE_P(out);
}

As you can see, in this skeleton we’re dealing with a data type where \( \text{union}(X, Y, Z) = \text{union}(\text{union}(X, Y), Z) \). It’s easy enough to support data types where this is not the case, by implementing the proper union algorithm in this GiST support method.

The result of the \text{union} function must be a value of the index’s storage type, whatever that is (it might or might not be different from the indexed column’s type). The \text{union} function should return a pointer to newly \text{palloc}ed memory. You can’t just return the input value as-is, even if there is no type change.

As shown above, the \text{union} function’s first \text{internal} argument is actually a GistEntryVector pointer. The second argument is a pointer to an integer variable, which can be ignored. (It used to be required that the \text{union} function store the size of its result value into that variable, but this is no longer necessary.)

\text{compress}

Converts the data item into a format suitable for physical storage in an index page.

The SQL declaration of the function must look like this:

CREATE OR REPLACE FUNCTION my_compress(internal)
RETURNS internal
Chapter 61. GiST Indexes

As 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:

PG_FUNCTION_INFO_V1(my_compress);

Datum
my_compress(PG_FUNCTION_ARGS)
{
    GISTENTRY *entry = (GISTENTRY *) PG_GETARG_POINTER(0);
    GISTENTRY *retval;

    if (entry->leafkey)
    {
        /* replace entry->key with a compressed version */
        compressed_data_type *compressed_data = palloc(sizeof(compressed_data_type));
        /* fill *compressed_data from entry->key ... */

        retval = palloc(sizeof(GISTENTRY));
        gistentryinit(*retval, PointerGetDatum(compressed_data),
                       entry->rel, entry->page, entry->offset, FALSE);
    }
    else
    {
        /* typically we needn’t do anything with non-leaf entries */
        retval = entry;
    }

    PG_RETURN_POINTER(retval);
}

You have to adapt compressed_data_type to the specific type you’re converting to in order to compress your leaf nodes, of course.

decompress

The reverse of the compress method. Converts the index representation of the data item into a format that can be manipulated by the other GiST methods in the operator class.

The SQL declaration of the function must look like this:

CREATE OR REPLACE FUNCTION my_decompress(internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;

And the matching code in the C module could then follow this skeleton:

PG_FUNCTION_INFO_V1(my_decompress);

Datum
my_decompress(PG_FUNCTION_ARGS)
{
    PG_RETURN_POINTER(PG_GETARG_POINTER(0));
}

The above skeleton is suitable for the case where no decompression is needed.
**Chapter 61. GiST Indexes**

penalty

Returns a value indicating the “cost” of inserting the new entry into a particular branch of the tree. Items will be inserted down the path of least penalty in the tree. Values returned by penalty should be non-negative. If a negative value is returned, it will be treated as zero.

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_penalty(internal, internal, internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;  -- in some cases penalty functions need not be strict
```

And the matching code in the C module could then follow this skeleton:

```c
Datum
my_penalty(PG_FUNCTION_ARGS)
{
    GISTENTRY *origentry = (GISTENTRY *) PG_GETARG_POINTER(0);
    GISTENTRY *newentry = (GISTENTRY *) PG_GETARG_POINTER(1);
    float penalty = (float *) PG_GETARG_POINTER(2);
    data_type *orig = DatumGetDataType(origentry->key);
    data_type *new = DatumGetDataType(newentry->key);

    *penalty = my_penalty_implementation(orig, new);
    PG_RETURN_POINTER(penalty);
}
```

For historical reasons, the penalty function doesn’t just return a float result; instead it has to store the value at the location indicated by the third argument. The return value per se is ignored, though it’s conventional to pass back the address of that argument.

The penalty function is crucial to good performance of the index. It’ll get used at insertion time to determine which branch to follow when choosing where to add the new entry in the tree. At query time, the more balanced the index, the quicker the lookup.

picksplit

When an index page split is necessary, this function decides which entries on the page are to stay on the old page, and which are to move to the new page.

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_picksplit(internal, internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;
```

And the matching code in the C module could then follow this skeleton:

```c
Datum
my_picksplit(PG_FUNCTION_ARGS)
{
    GistEntryVector *entryvec = (GistEntryVector *) PG_GETARG_POINTER(0);
    GIST_SPLITVEC *v = (GIST_SPLITVEC *) PG_GETARG_POINTER(1);
    OffsetNumber maxoff = entryvec->n - 1;
    GISTENTRY *ent = entryvec->vector;
    int i,
        nbytes;
    OffsetNumber *left,
```
*right;
data_type  *tmp_union;
data_type  *unionL;
data_type  *unionR;
GISTENTRY **raw_entryvec;

maxoff = entryvec->n - 1;
nbytes = (maxoff + 1) * sizeof(OffsetNumber);

v->spl_left = (OffsetNumber *) palloc(nbytes);
left = v->spl_left;
v->spl_nleft = 0;

v->spl_right = (OffsetNumber *) palloc(nbytes);
right = v->spl_right;
v->spl_nright = 0;

unionL = NULL;
unionR = NULL;

/* Initialize the raw entry vector. */
raw_entryvec = (GISTENTRY **) malloc(entryvec->n * sizeof(void *));
for (i = FirstOffsetNumber; i <= maxoff; i = OffsetNumberNext(i))
    raw_entryvec[i] = &(entryvec->vector[i]);

for (i = FirstOffsetNumber; i <= maxoff; i = OffsetNumberNext(i))
{
    int real_index = raw_entryvec[i] - entryvec->vector;

    tmp_union = DatumGetDataType(entryvec->vector[real_index].key);
    Assert(tmp_union != NULL);

    /*
     * Choose where to put the index entries and update unionL and unionR
     * accordingly. Append the entries to either v->spl_left or
     * v->spl_right, and care about the counters.
     */
    if (my_choice_is_left(unionL, curl, unionR, curr))
    {
        if (unionL == NULL)
            unionL = tmp_union;
        else
            unionL = my_union_implementation(unionL, tmp_union);

        *left = real_index;
        ++left;
        ++(v->spl_nleft);
    }
    else
    {
        /*
         * Same on the right
         */
    }
}
v->spl_ldatum = DataTypeGetDatum(unionL);
v->spl_rdatum = DataTypeGetDatum(unionR);
PG_RETURN_POINTER(v);
}

Notice that the `picksplit` function’s result is delivered by modifying the passed-in v structure. The return value per se is ignored, though it’s conventional to pass back the address of v.

Like `penalty`, the `picksplit` function is crucial to good performance of the index. Designing suitable `penalty` and `picksplit` implementations is where the challenge of implementing well-performing GiST indexes lies.

`same`

Returns true if two index entries are identical, false otherwise. (An “index entry” is a value of the index’s storage type, not necessarily the original indexed column’s type.)

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_same(storage_type, storage_type, internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;
```

And the matching code in the C module could then follow this skeleton:

```c
PG_FUNCTION_INFO_V1(my_same);

Datum
my_same(PG_FUNCTION_ARGS)
{
    prefix_range *v1 = PG_GETARG_PREFIX_RANGE_P(0);
    prefix_range *v2 = PG_GETARG_PREFIX_RANGE_P(1);
    bool *result = (bool *) PG_GETARG_POINTER(2);

    *result = my_eq(v1, v2);
    PG_RETURN_POINTER(result);
}
```

For historical reasons, the `same` function doesn’t just return a Boolean result; instead it has to store the flag at the location indicated by the third argument. The return value per se is ignored, though it’s conventional to pass back the address of that argument.

`distance`

Given an index entry p and a query value q, this function determines the index entry’s “distance” from the query value. This function must be supplied if the operator class contains any ordering operators. A query using the ordering operator will be implemented by returning index entries with the smallest “distance” values first, so the results must be consistent with the operator’s semantics. For a leaf index entry the result just represents the distance to the index entry; for an internal tree node, the result must be the smallest distance that any child entry could have.

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_distance(internal, data_type, smallint, oid, internal)
RETURNS float8
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;
```

And the matching code in the C module could then follow this skeleton:

```c
PG_FUNCTION_INFO_V1(my_distance);

Datum
my_distance(PG_FUNCTION_ARGS)
```
The arguments to the distance function are identical to the arguments of the consistent function.

Some approximation is allowed when determining the distance, so long as the result is never greater than the entry’s actual distance. Thus, for example, distance to a bounding box is usually sufficient in geometric applications. For an internal tree node, the distance returned must not be greater than the distance to any of the child nodes. If the returned distance is not exact, the function must set *recheck to true. (This is not necessary for internal tree nodes; for them, the calculation is always assumed to be inexact.) In this case the executor will calculate the accurate distance after fetching the tuple from the heap, and reorder the tuples if necessary.

If the distance function returns *recheck = true for any leaf node, the original ordering operator’s return type must be float8 or float4, and the distance function’s result values must be comparable to those of the original ordering operator, since the executor will sort using both distance function results and recalculated ordering-operator results. Otherwise, the distance function’s result values can be any finite float8 values, so long as the relative order of the result values matches the order returned by the ordering operator. (Infinity and minus infinity are used internally to handle cases such as nulls, so it is not recommended that distance functions return these values.)

The matching code in the C module could then follow this skeleton:

```c
PG_FUNCTION_INFO_V1(my_fetch);

Datum
my_fetch(PG_FUNCTION_ARGS)
{
   // Implementation...
}
```

fetch

Converts the compressed index representation of a data item into the original data type, for index-only scans. The returned data must be an exact, non-lossy copy of the originally indexed value.

The SQL declaration of the function must look like this:

```sql
CREATE OR REPLACE FUNCTION my_fetch(internal)
RETURNS internal
AS 'MODULE_PATHNAME'
LANGUAGE C STRICT;
```

The argument is a pointer to a GISTENTRY struct. On entry, its key field contains a non-NULL leaf datum in compressed form. The return value is another GISTENTRY struct, whose key field contains the same datum in its original, uncompressed form. If the opclass’s compress function does nothing for leaf entries, the fetch method can return the argument as-is.

The matching code in the C module could then follow this skeleton:

```c
PG_FUNCTION_INFO_V1(my_fetch);

Datum
my_fetch(PG_FUNCTION_ARGS)
{
   // Implementation...
}
```
GISTENTRY *entry = (GISTENTRY *) PG_GETARG_POINTER(0);
input_data_type *in = DatumGetPointer(entry->key);
fetched_data_type *fetched_data;
GISTENTRY *retval;

retval = palloc(sizeof(GISTENTRY));
fetched_data = palloc(sizeof(fetched_data_type));

/* Convert 'fetched_data' into the a Datum of the original datatype. */
/* fill *retval from fetched_data. */
gistentryinit(*retval, PointerGetDatum(converted_datum),
    entry->rel, entry->page, entry->offset, FALSE);

PG_RETURN_POINTER(retval);
}

If the compress method is lossy for leaf entries, the operator class cannot support index-only scans, and must not define a fetch function.

All the GiST support methods are normally called in short-lived memory contexts; that is, CurrentMemoryContext will get reset after each tuple is processed. It is therefore not very important to worry about pfree’ing everything you palloc. However, in some cases it’s useful for a support method to cache data across repeated calls. To do that, allocate the longer-lived data in fcinfo->flinfo->fn_mcxt, and keep a pointer to it in fcinfo->flinfo->fn_extra. Such data will survive for the life of the index operation (e.g., a single GiST index scan, index build, or index tuple insertion). Be careful to pfree the previous value when replacing a fn_extra value, or the leak will accumulate for the duration of the operation.

61.4. Implementation

61.4.1. GiST buffering build

Building large GiST indexes by simply inserting all the tuples tends to be slow, because if the index tuples are scattered across the index and the index is large enough to not fit in cache, the insertions need to perform a lot of random I/O. Beginning in version 9.2, PostgreSQL supports a more efficient method to build GiST indexes based on buffering, which can dramatically reduce the number of random I/Os needed for non-ordered data sets. For well-ordered data sets the benefit is smaller or non-existent, because only a small number of pages receive new tuples at a time, and those pages fit in cache even if the index as whole does not.

However, buffering index build needs to call the penalty function more often, which consumes some extra CPU resources. Also, the buffers used in the buffering build need temporary disk space, up to the size of the resulting index. Buffering can also influence the quality of the resulting index, in both positive and negative directions. That influence depends on various factors, like the distribution of the input data and the operator class implementation.

By default, a GiST index build switches to the buffering method when the index size reaches effective_cache_size. It can be manually turned on or off by the buffering parameter to the CREATE INDEX command. The default behavior is good for most cases, but turning buffering off might speed up the build somewhat if the input data is ordered.
61.5. Examples

The PostgreSQL source distribution includes several examples of index methods implemented using GiST. The core system currently provides text search support (indexing for `tsvector` and `tsquery`) as well as R-Tree equivalent functionality for some of the built-in geometric data types (see `src/backend/access/gist/gistproc.c`). The following contrib modules also contain GiST operator classes:

- `btree_gist`  
  B-tree equivalent functionality for several data types
- `cube`  
  Indexing for multidimensional cubes
- `hstore`  
  Module for storing (key, value) pairs
- `intarray`  
  RD-Tree for one-dimensional array of int4 values
- `ltree`  
  Indexing for tree-like structures
- `pg_trgm`  
  Text similarity using trigram matching
- `seg`  
  Indexing for “float ranges”
Chapter 62. SP-GiST Indexes

62.1. Introduction

SP-GiST is an abbreviation for space-partitioned GiST. SP-GiST supports partitioned search trees, which facilitate development of a wide range of different non-balanced data structures, such as quad-trees, k-d trees, and radix trees (tries). The common feature of these structures is that they repeatedly divide the search space into partitions that need not be of equal size. Searches that are well matched to the partitioning rule can be very fast.

These popular data structures were originally developed for in-memory usage. In main memory, they are usually designed as a set of dynamically allocated nodes linked by pointers. This is not suitable for direct storing on disk, since these chains of pointers can be rather long which would require too many disk accesses. In contrast, disk-based data structures should have a high fanout to minimize I/O. The challenge addressed by SP-GiST is to map search tree nodes to disk pages in such a way that a search need access only a few disk pages, even if it traverses many nodes.

Like GiST, SP-GiST is meant to allow the development of custom data types with the appropriate access methods, by an expert in the domain of the data type, rather than a database expert.

Some of the information here is derived from Purdue University’s SP-GiST Indexing Project web site¹. The SP-GIST implementation in PostgreSQL is primarily maintained by Teodor Sigaev and Oleg Bartunov, and there is more information on their web site².

62.2. Built-in Operator Classes

The core PostgreSQL distribution includes the SP-GiST operator classes shown in Table 62-1.

Table 62-1. Built-in SP-GiST Operator Classes

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>kd_point_ops</td>
<td>point</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>quad_point_ops</td>
<td>point</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>range_ops</td>
<td>any range type</td>
<td>&amp;&amp;</td>
</tr>
<tr>
<td>box_ops</td>
<td>box</td>
<td>&lt;&lt;</td>
</tr>
<tr>
<td>text_ops</td>
<td>text</td>
<td>&lt;=</td>
</tr>
</tbody>
</table>

Of the two operator classes for type point, quad_point_ops is the default. kd_point_ops supports the same operators but uses a different index data structure that may offer better performance in some applications.

62.3. Extensibility

SP-GiST offers an interface with a high level of abstraction, requiring the access method developer to implement only methods specific to a given data type. The SP-GiST core is responsible for efficient disk mapping and searching the tree structure. It also takes care of concurrency and logging considerations.

Leaf tuples of an SP-GiST tree contain values of the same data type as the indexed column. Leaf tuples at the root level will always contain the original indexed data value, but leaf tuples at lower levels might contain only a compressed representation, such as a suffix. In that case the operator class support functions must be able to reconstruct the original value using information accumulated from the inner tuples that are passed through to reach the leaf level.

Inner tuples are more complex, since they are branching points in the search tree. Each inner tuple contains a set of one or more nodes, which represent groups of similar leaf values. A node contains a downlink that leads either to another, lower-level inner tuple, or to a short list of leaf tuples that all lie on the same index page. Each node normally has a label that describes it; for example, in a radix tree the node label could be the next character of the string value. (Alternatively, an operator class can omit the node labels, if it works with a fixed set of nodes for all inner tuples; see Section 62.4.2.) Optionally, an inner tuple can have a prefix value that describes all its members. In a radix tree this could be the common prefix of the represented strings. The prefix value is not necessarily really a prefix, but can be any data needed by the operator class; for example, in a quad-tree it can store the central point that the four quadrants are measured with respect to. A quad-tree inner tuple would then also contain four nodes corresponding to the quadrants around this central point.

Some tree algorithms require knowledge of level (or depth) of the current tuple, so the SP-GiST core provides the possibility for operator classes to manage level counting while descending the tree. There is also support for incrementally reconstructing the represented value when that is needed, and for passing down additional data (called traverse values) during a tree descent.

Note: The SP-GiST core code takes care of null entries. Although SP-GiST indexes do store entries for nulls in indexed columns, this is hidden from the index operator class code: no null index entries or search conditions will ever be passed to the operator class methods. (It is assumed that SP-GiST operators are strict and so cannot succeed for null values.) Null values are therefore not discussed further here.

There are five user-defined methods that an index operator class for SP-GiST must provide. All five follow the convention of accepting two internal arguments, the first of which is a pointer to a C struct containing input values for the support method, while the second argument is a pointer to a C struct where output values must be placed. Four of the methods just return void, since all their results appear in the output struct; but leaf_consistent additionally returns a boolean result. The methods must not modify any fields of their input structs. In all cases, the output struct is initialized to zeroes before calling the user-defined method.

The five user-defined methods are:

**config**

Returns static information about the index implementation, including the data type OIDs of the prefix and node label data types.

The SQL declaration of the function must look like this:

```
CREATE FUNCTION my_config(internal, internal) RETURNS void ...
```
Chapter 62. SP-GiST Indexes

The first argument is a pointer to a `spgConfigIn` C struct, containing input data for the function. The second argument is a pointer to a `spgConfigOut` C struct, which the function must fill with result data.

```c
typedef struct spgConfigIn {
    Oid attType; /* Data type to be indexed */
} spgConfigIn;

typedef struct spgConfigOut {
    Oid prefixType; /* Data type of inner-tuple prefixes */
    Oid labelType; /* Data type of inner-tuple node labels */
    bool canReturnData; /* Opclass can reconstruct original data */
    bool longValuesOK; /* Opclass can cope with values > 1 page */
} spgConfigOut;
```

`attType` is passed in order to support polymorphic index operator classes; for ordinary fixed-data-type operator classes, it will always have the same value and so can be ignored.

For operator classes that do not use prefixes, `prefixType` can be set to `VOIDOID`. Likewise, for operator classes that do not use node labels, `labelType` can be set to `VOIDOID`. `canReturnData` should be set true if the operator class is capable of reconstructing the originally-supplied index value. `longValuesOK` should be set true only when the `attType` is of variable length and the operator class is capable of segmenting long values by repeated suffixing (see Section 62.4.1).

`choose`

Chooses a method for inserting a new value into an inner tuple.

The SQL declaration of the function must look like this:

```sql
CREATE FUNCTION my_choose(internal, internal) RETURNS void ...
```

The first argument is a pointer to a `spgChooseIn` C struct, containing input data for the function. The second argument is a pointer to a `spgChooseOut` C struct, which the function must fill with result data.

```c
typedef struct spgChooseIn {
    Datum datum; /* original datum to be indexed */
    Datum leafDatum; /* current datum to be stored at leaf */
    int level; /* current level (counting from zero) */
    /* Data from current inner tuple */
    bool allTheSame; /* tuple is marked all-the-same? */
    bool hasPrefix; /* tuple has a prefix? */
    Datum prefixDatum; /* if so, the prefix value */
    int nNodes; /* number of nodes in the inner tuple */
    Datum *nodeLabels; /* node label values (NULL if none) */
} spgChooseIn;

typedef enum spgChooseResultType {
    spgMatchNode = 1, /* descend into existing node */
    spgAddNode, /* add a node to the inner tuple */
    spgSplitTuple /* split inner tuple (change its prefix) */
} spgChooseResultType;

typedef struct spgChooseOut
Chapter 62. SP-GiST Indexes

```c
{
    spgChooseResultType resultType; /* action code, see above */
    union
    {
        struct /* results for spgMatchNode */
        {
            int nodeN; /* descend to this node (index from 0) */
            int levelAdd; /* increment level by this much */
            Datum restDatum; /* new leaf datum */
        } matchNode;
        struct /* results for spgAddNode */
        {
            Datum nodeLabel; /* new node’s label */
            int nodeN; /* where to insert it (index from 0) */
        } addNode;
        struct /* results for spgSplitTuple */
        {
            /* Info to form new inner tuple with one node */
            bool prefixHasPrefix; /* tuple should have a prefix? */
            Datum prefixPrefixDatum; /* if so, its value */
            Datum nodeLabel; /* node’s label */
            /* Info to form new lower-level inner tuple with all old nodes */
            bool postfixHasPrefix; /* tuple should have a prefix? */
            Datum postfixPrefixDatum; /* if so, its value */
        } splitTuple;
    } result;
} spgChooseOut;
```

datum is the original datum that was to be inserted into the index. leafDatum is initially the same as datum, but can change at lower levels of the tree if the choose or picksplit methods change it. When the insertion search reaches a leaf page, the current value of leafDatum is what will be stored in the newly created leaf tuple. level is the current inner tuple’s level, starting at zero for the root level. allTheSame is true if the current inner tuple is marked as containing multiple equivalent nodes (see Section 62.4.3). hasPrefix is true if the current inner tuple contains a prefix; if so, prefixDatum is its value. nNodes is the number of child nodes contained in the inner tuple, and nodeLabels is an array of their label values, or NULL if there are no labels.

The choose function can determine either that the new value matches one of the existing child nodes, or that a new child node must be added, or that the new value is inconsistent with the tuple prefix and so the inner tuple must be split to create a less restrictive prefix.

If the new value matches one of the existing child nodes, set resultType to spgMatchNode. Set nodeN to the index (from zero) of that node in the node array. Set levelAdd to the increment in level caused by descending through that node, or leave it as zero if the operator class does not use levels. Set restDatum to equal datum if the operator class does not modify datums from one level to the next, or otherwise set it to the modified value to be used as leafDatum at the next level.

If a new child node must be added, set resultType to spgAddNode. Set nodeLabel to the label to be used for the new node, and set nodeN to the index (from zero) at which to insert the node in the node array. After the node has been added, the choose function will be called again with the modified inner tuple; that call should result in an spgMatchNode result.

If the new value is inconsistent with the tuple prefix, set resultType to spgSplitTuple. This action moves all the existing nodes into a new lower-level inner tuple, and replaces the existing inner tuple with a tuple having a single node that links to the new lower-level inner tuple. Set
prefixHasPrefix to indicate whether the new upper tuple should have a prefix, and if so set prefixPrefixDatum to the prefix value. This new prefix value must be sufficiently less restrictive than the original to accept the new value to be indexed, and it should be no longer than the original prefix. Set nodeLabel to the label to be used for the node that will point to the new lower-level inner tuple. Set postfixHasPrefix to indicate whether the new lower-level inner tuple should have a prefix, and if so set postfixPrefixDatum to the prefix value. The combination of these two prefixes and the additional label must have the same meaning as the original prefix, because there is no opportunity to alter the node labels that are moved to the new lower-level tuple, nor to change any child index entries. After the node has been split, the choose function will be called again with the replacement inner tuple. That call will usually result in an spgAddNode result, since presumably the node label added in the split step will not match the new value; so after that, there will be a third call that finally returns spgMatchNode and allows the insertion to descend to the leaf level.

picksplit

Decides how to create a new inner tuple over a set of leaf tuples.

The SQL declaration of the function must look like this:

```sql
CREATE FUNCTION my_picksplit(internal, internal) RETURNS void ... 
```

The first argument is a pointer to a spgPickSplitIn C struct, containing input data for the function. The second argument is a pointer to a spgPickSplitOut C struct, which the function must fill with result data.

```c
typedef struct spgPickSplitIn
{
    int nTuples; /* number of leaf tuples */
    Datum *datums; /* their datums (array of length nTuples) */
    int level; /* current level (counting from zero) */
} spgPickSplitIn;

typedef struct spgPickSplitOut
{
    bool hasPrefix; /* new inner tuple should have a prefix? */
    Datum prefixDatum; /* if so, its value */
    int nNodes; /* number of nodes for new inner tuple */
    Datum *nodeLabels; /* their labels (or NULL for no labels) */
    int *mapTuplesToNodes; /* node index for each leaf tuple */
    Datum *leafTupleDatums; /* datum to store in each new leaf tuple */
} spgPickSplitOut;
```

nTuples is the number of leaf tuples provided. datums is an array of their datum values. level is the current level that all the leaf tuples share, which will become the level of the new inner tuple.

Set hasPrefix to indicate whether the new inner tuple should have a prefix, and if so set prefixDatum to the prefix value. Set nNodes to indicate the number of nodes that the new inner tuple will contain, and set nodeLabels to an array of their label values, or to NULL if node labels are not required. Set mapTuplesToNodes to an array that gives the index (from zero) of the node that each leaf tuple should be assigned to. Set leafTupleDatums to an array of the values to be stored in the new leaf tuples (these will be the same as the input datums if the operator class does not modify datums from one level to the next). Note that the picksplit function is responsible for palloc'ing the nodeLabels, mapTuplesToNodes and leafTupleDatums arrays.
If more than one leaf tuple is supplied, it is expected that the \texttt{picksplit} function will classify them into more than one node; otherwise it is not possible to split the leaf tuples across multiple pages, which is the ultimate purpose of this operation. Therefore, if the \texttt{picksplit} function ends up placing all the leaf tuples in the same node, the core SP-GiST code will override that decision and generate an inner tuple in which the leaf tuples are assigned at random to several identically-labeled nodes. Such a tuple is marked \texttt{allTheSame} to signify that this has happened. The \texttt{choose} and \texttt{inner\_consistent} functions must take suitable care with such inner tuples. See Section 62.4.3 for more information.

\texttt{picksplit} can be applied to a single leaf tuple only in the case that the \texttt{config} function set \texttt{longValuesOK} to true and a larger-than-a-page input value has been supplied. In this case the point of the operation is to strip off a prefix and produce a new, shorter leaf datum value. The call will be repeated until a leaf datum short enough to fit on a page has been produced. See Section 62.4.1 for more information.

\texttt{inner\_consistent}

Returns set of nodes (branches) to follow during tree search.

The SQL declaration of the function must look like this:

\begin{verbatim}
CREATE FUNCTION my\_inner\_consistent(internal, internal) RETURNS void ...
\end{verbatim}

The first argument is a pointer to a \texttt{spgInnerConsistentIn} C struct, containing input data for the function. The second argument is a pointer to a \texttt{spgInnerConsistentOut} C struct, which the function must fill with result data.

\begin{verbatim}
typedef struct spgInnerConsistentIn
{
  ScanKey scankeys; /* array of operators and comparison values */
  int nkeys; /* length of array */
  Datum reconstructedValue; /* value reconstructed at parent */
  void *traversalValue; /* opclass-specific traverse value */
  MemoryContext traversalMemoryContext; /* put new traverse values here */
  int level; /* current level (counting from zero) */
  bool returnData; /* original data must be returned? */
  /* Data from current inner tuple */
  bool allTheSame; /* tuple is marked all-the-same? */
  bool hasPrefix; /* tuple has a prefix? */
  Datum prefixDatum; /* if so, the prefix value */
  int nNodes; /* number of nodes in the inner tuple */
  Datum *nodeLabels; /* node label values (NULL if none) */
} spgInnerConsistentIn;

typedef struct spgInnerConsistentOut
{
  int nNodes; /* number of child nodes to be visited */
  int *nodeNumbers; /* their indexes in the node array */
  int *levelAdds; /* increment level by this much for each */
  Datum *reconstructedValues; /* associated reconstructed values */
  void **traversalValues; /* opclass-specific traverse values */
} spgInnerConsistentOut;
\end{verbatim}

The array \texttt{scankeys}, of length \texttt{nkeys}, describes the index search condition(s). These conditions are combined with AND — only index entries that satisfy all of them are interesting. (Note that \texttt{nkeys = 0} implies that all index entries satisfy the query.) Usually the consistent function only cares about the \texttt{sk\_strategy} and \texttt{sk\_argument} fields of each array entry, which respectively give the indexable operator and comparison value. In particular it is not necessary to check
sk_flags to see if the comparison value is NULL, because the SP-GiST core code will filter out such conditions. reconstructedValue is the value reconstructed for the parent tuple; it is (Datum) 0 at the root level or if the inner_consistent function did not provide a value at the parent level. traversalValue is a pointer to any traverse data passed down from the previous call of inner_consistent on the parent index tuple, or NULL at the root level. traversalMemoryContext is the memory context in which to store output traverse values (see below). level is the current inner tuple’s level, starting at zero for the root level. returnData is true if reconstructed data is required for this query; this will only be so if the config function asserted canReturnData. allTheSame is true if the current inner tuple is marked “all-the-same”; in this case all the nodes have the same label (if any) and so either all or none of them match the query (see Section 62.4.3). hasPrefix is true if the current inner tuple contains a prefix; if so, prefixDatum is its value. nNodes is the number of child nodes contained in the inner tuple, and nodeLabels is an array of their label values, or NULL if the nodes do not have labels.

nNodes must be set to the number of child nodes that need to be visited by the search, and nodeNumbers must be set to an array of their indexes. If the operator class keeps track of levels, set levelAdds to an array of the level increments required when descending to each node to be visited. (Often these increments will be the same for all the nodes, but that’s not necessarily so, so an array is used.) If value reconstruction is needed, set reconstructedValues to an array of the values reconstructed for each child node to be visited; otherwise, leave reconstructedValues as NULL. If it is desired to pass down additional out-of-band information (“traverse values”) to lower levels of the tree search, set traversalValues to an array of the appropriate traverse values, one for each child node to be visited; otherwise, leave traversalValues as NULL. Note that the inner_consistent function is responsible for palloc’ing the nodeNumbers, levelAdds, reconstructedValues, and traversalValues arrays in the current memory context. However, any output traverse values pointed to by the traversalValues array should be allocated in traversalMemoryContext. Each traverse value must be a single palloc’d chunk.

leaf_consistent

Returns true if a leaf tuple satisfies a query.

The SQL declaration of the function must look like this:

```sql
CREATE FUNCTION my_leaf_consistent(internal, internal) RETURNS bool ...
```

The first argument is a pointer to a spgLeafConsistentIn C struct, containing input data for the function. The second argument is a pointer to a spgLeafConsistentOut C struct, which the function must fill with result data.

typedef struct spgLeafConsistentIn
{
    ScanKey scankeys;         /* array of operators and comparison values */
    int nkeys;                /* length of array */
    void *traversalValue;     /* opclass-specific traverse value */
    Datum reconstructedValue; /* value reconstructed at parent */
    int level;                /* current level (counting from zero) */
    bool returnData;          /* original data must be returned? */
    Datum leafDatum;          /* datum in leaf tuple */
} spgLeafConsistentIn;

typedef struct spgLeafConsistentOut
{
    Datum leafValue;          /* reconstructed original data, if any */
} spgLeafConsistentOut;
bool recheck; /* set true if operator must be rechecked */
}

The array scankeys, of length nkeys, describes the index search condition(s). These conditions are combined with AND — only index entries that satisfy all of them satisfy the query. (Note that nkeys = 0 implies that all index entries satisfy the query.) Usually the consistent function only cares about the sk_strategy and sk_argument fields of each array entry, which respectively give the indexable operator and comparison value. In particular it is not necessary to check sk_flags to see if the comparison value is NULL, because the SP-GiST core code will filter out such conditions. reconstructedValue is the value reconstructed for the parent tuple; it is (Datum) 0 at the root level or if the inner consistent function did not provide a value at the parent level. traversalValue is a pointer to any traverse data passed down from the previous call of inner consistent on the parent index tuple, or NULL at the root level. level is the current leaf tuple’s level, starting at zero for the root level. returnData is true if reconstructed data is required for this query; this will only be so if the config function asserted canReturnData. leafDatum is the key value stored in the current leaf tuple.

The function must return true if the leaf tuple matches the query, or false if not. In the true case, if returnData is true then leafValue must be set to the value originally supplied to be indexed for this leaf tuple. Also, recheck may be set to true if the match is uncertain and so the operator(s) must be re-applied to the actual heap tuple to verify the match.

All the SP-GiST support methods are normally called in a short-lived memory context; that is, CurrentMemoryContext will be reset after processing of each tuple. It is therefore not very important to worry about free’ing everything you malloc. (The config method is an exception: it should try to avoid leaking memory. But usually the config method need do nothing but assign constants into the passed parameter struct.)

If the indexed column is of a collatable data type, the index collation will be passed to all the support methods, using the standard PG_GET_COLLATION() mechanism.

### 62.4. Implementation

This section covers implementation details and other tricks that are useful for implementers of SP-GiST operator classes to know.

#### 62.4.1. SP-GiST Limits

Individual leaf tuples and inner tuples must fit on a single index page (8kB by default). Therefore, when indexing values of variable-length data types, long values can only be supported by methods such as radix trees, in which each level of the tree includes a prefix that is short enough to fit on a page, and the final leaf level includes a suffix also short enough to fit on a page. The operator class should set longValuesOK to TRUE only if it is prepared to arrange for this to happen. Otherwise, the SP-GiST core will reject any request to index a value that is too large to fit on an index page.

Likewise, it is the operator class’s responsibility that inner tuples do not grow too large to fit on an index page; this limits the number of child nodes that can be used in one inner tuple, as well as the maximum size of a prefix value.

Another limitation is that when an inner tuple’s node points to a set of leaf tuples, those tuples must all be in the same index page. (This is a design decision to reduce seeking and save space in the links that chain such tuples together.) If the set of leaf tuples grows too large for a page, a split is performed and an intermediate inner tuple is inserted. For this to fix the problem, the new inner tuple must divide
Chapter 62. SP-GiST Indexes

the set of leaf values into more than one node group. If the operator class’s \texttt{picksplit} function fails to do that, the SP-GiST core resorts to extraordinary measures described in Section 62.4.3.

### 62.4.2. SP-GiST Without Node Labels

Some tree algorithms use a fixed set of nodes for each inner tuple; for example, in a quad-tree there are always exactly four nodes corresponding to the four quadrants around the inner tuple’s centroid point. In such a case the code typically works with the nodes by number, and there is no need for explicit node labels. To suppress node labels (and thereby save some space), the \texttt{picksplit} function can return NULL for the \texttt{nodeLabels} array. This will in turn result in \texttt{nodeLabels} being NULL during subsequent calls to \texttt{choose} and \texttt{inner\_consistent}. In principle, node labels could be used for some inner tuples and omitted for others in the same index.

When working with an inner tuple having unlabeled nodes, it is an error for \texttt{choose} to return \texttt{spgAddNode}, since the set of nodes is supposed to be fixed in such cases. Also, there is no provision for generating an unlabeled node in \texttt{spgSplitTuple} actions, since it is expected that an \texttt{spgAddNode} action will be needed as well.

### 62.4.3. “All-the-same” Inner Tuples

The SP-GiST core can override the results of the operator class’s \texttt{picksplit} function when \texttt{picksplit} fails to divide the supplied leaf values into at least two node categories. When this happens, the new inner tuple is created with multiple nodes that each have the same label (if any) that \texttt{picksplit} gave to the one node it did use, and the leaf values are divided at random among these equivalent nodes. The \texttt{allTheSame} flag is set on the inner tuple to warn the \texttt{choose} and \texttt{inner\_consistent} functions that the tuple does not have the node set that they might otherwise expect.

When dealing with an \texttt{allTheSame} tuple, a \texttt{choose} result of \texttt{spgMatchNode} is interpreted to mean that the new value can be assigned to any of the equivalent nodes; the core code will ignore the supplied \texttt{nodeN} value and descend into one of the nodes at random (so as to keep the tree balanced). It is an error for \texttt{choose} to return \texttt{spgAddNode}, since that would make the nodes not all equivalent; the \texttt{spgSplitTuple} action must be used if the value to be inserted doesn’t match the existing nodes.

When dealing with an \texttt{allTheSame} tuple, the \texttt{inner\_consistent} function should return either all or none of the nodes as targets for continuing the index search, since they are all equivalent. This may or may not require any special-case code, depending on how much the \texttt{inner\_consistent} function normally assumes about the meaning of the nodes.

### 62.5. Examples

The PostgreSQL source distribution includes several examples of index operator classes for SP-GiST, as described in Table 62-1. Look into \texttt{src/backend/access/spgist/} and \texttt{src/backend/utils/adt/} to see the code.
### Chapter 63. GIN Indexes

#### 63.1. Introduction

GIN stands for Generalized Inverted Index. GIN is designed for handling cases where the items to be indexed are composite values, and the queries to be handled by the index need to search for element values that appear within the composite items. For example, the items could be documents, and the queries could be searches for documents containing specific words.

We use the word *item* to refer to a composite value that is to be indexed, and the word *key* to refer to an element value. GIN always stores and searches for keys, not item values per se.

A GIN index stores a set of (key, posting list) pairs, where a *posting list* is a set of row IDs in which the key occurs. The same row ID can appear in multiple posting lists, since an item can contain more than one key. Each key value is stored only once, so a GIN index is very compact for cases where the same key appears many times.

GIN is generalized in the sense that the GIN access method code does not need to know the specific operations that it accelerates. Instead, it uses custom strategies defined for particular data types. The strategy defines how keys are extracted from indexed items and query conditions, and how to determine whether a row that contains some of the key values in a query actually satisfies the query.

One advantage of GIN is that it allows the development of custom data types with the appropriate access methods, by an expert in the domain of the data type, rather than a database expert. This is much the same advantage as using GiST.

The GIN implementation in PostgreSQL is primarily maintained by Teodor Sigaev and Oleg Bartunov. There is more information about GIN on their website[^1].

#### 63.2. Built-in Operator Classes

The core PostgreSQL distribution includes the GIN operator classes shown in Table 63-1. (Some of the optional modules described in Appendix F provide additional GIN operator classes.)

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>_abstime_ops</td>
<td>abstime[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
<tr>
<td>_bit_ops</td>
<td>bit[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
<tr>
<td>_bool_ops</td>
<td>boolean[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
<tr>
<td>_bpchar_ops</td>
<td>character[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
<tr>
<td>_bytea_ops</td>
<td>bytea[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
<tr>
<td>_char_ops</td>
<td>&quot;char&quot;[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
<tr>
<td>_cidr_ops</td>
<td>cidr[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
<tr>
<td>_date_ops</td>
<td>date[]</td>
<td>&amp;&amp;, &lt;@ - @&gt;</td>
</tr>
</tbody>
</table>

[^1]: http://www.sai.msu.su/~megera/wiki/Gin
### Chapter 63. GIN Indexes

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>_float4_ops</td>
<td>float4[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_float8_ops</td>
<td>float8[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_inet_ops</td>
<td>inet[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_int2_ops</td>
<td>smallint[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_int4_ops</td>
<td>integer[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_int8_ops</td>
<td>bigint[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_interval_ops</td>
<td>interval[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_macaddr_ops</td>
<td>macaddr[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_money_ops</td>
<td>money[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_name_ops</td>
<td>name[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_numeric_ops</td>
<td>numeric[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_oid_ops</td>
<td>oid[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_oidvector_ops</td>
<td>oidvector[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_reltime_ops</td>
<td>reltime[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_text_ops</td>
<td>text[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_time_ops</td>
<td>time[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_timestamp_ops</td>
<td>timestamp[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_timestamptz_ops</td>
<td>timestamp with time zone[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_timestz_ops</td>
<td>time with time zone[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_tinterval_ops</td>
<td>tinterval[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_varbit_ops</td>
<td>bit varying[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>_varchar_ops</td>
<td>character varying[]</td>
<td>&amp; &amp; &lt; @ = 0 &gt;</td>
</tr>
<tr>
<td>jsonb_ops</td>
<td>jsonb</td>
<td>@</td>
</tr>
<tr>
<td>jsonb_path_ops</td>
<td>jsonb</td>
<td>0 &gt;</td>
</tr>
<tr>
<td>tsvector_ops</td>
<td>tsvector</td>
<td>0 @ 0 0 0</td>
</tr>
</tbody>
</table>

Of the two operator classes for type `jsonb`, `jsonb_ops` is the default. `jsonb_path_ops` supports fewer operators but offers better performance for those operators. See Section 8.14.4 for details.

### 63.3. Extensibility

The GIN interface has a high level of abstraction, requiring the access method implementer only to implement the semantics of the data type being accessed. The GIN layer itself takes care of concurrency, logging and searching the tree structure.

All it takes to get a GIN access method working is to implement a few user-defined methods, which define the behavior of keys in the tree and the relationships between keys, indexed items, and indexable queries. In short, GIN combines extensibility with generality, code reuse, and a clean interface.

There are three methods that an operator class for GIN must provide:

```c
int compare(Datum a, Datum b)
```

Compares two keys (not indexed items!) and returns an integer less than zero, zero, or greater than zero, indicating whether the first key is less than, equal to, or greater than the second. Null
Chapter 63. GIN Indexes

keys are never passed to this function.

Datum *extractValue(Datum itemValue, int32 *nkeys, bool **nullFlags)

Returns a palloc’d array of keys given an item to be indexed. The number of returned keys must be stored into *nkeys. If any of the keys can be null, also palloc an array of *nkeys bool fields, store its address at *nullFlags, and set these null flags as needed. *nullFlags can be left NULL (its initial value) if all keys are non-null. The return value can be NULL if the item contains no keys.

Datum *extractQuery(Datum query, int32 *nkeys, StrategyNumber n, bool **pmatch, Pointer **extra_data, bool **nullFlags, int32 *searchMode)

Returns a palloc’d array of keys given a value to be queried; that is, query is the value on the right-hand side of an indexable operator whose left-hand side is the indexed column. n is the strategy number of the operator within the operator class (see Section 36.14.2). Often, extractQuery will need to consult n to determine the data type of query and the method it should use to extract key values. The number of returned keys must be stored into *nkeys. If any of the keys can be null, also palloc an array of *nkeys bool fields, store its address at *nullFlags, and set these null flags as needed. *nullFlags can be left NULL (its initial value) if all keys are non-null. The return value can be NULL if the query contains no keys.

searchMode is an output argument that allows extractQuery to specify details about how the search will be done. If *searchMode is set to GIN_SEARCH_MODE_DEFAULT (which is the value it is initialized to before call), only items that match at least one of the returned keys are considered candidate matches. If *searchMode is set to GIN_SEARCH_MODE_INCLUDE_EMPTY, then in addition to items containing at least one matching key, items that contain no keys at all are considered candidate matches. (This mode is useful for implementing is-subset-of operators, for example.) If *searchMode is set to GIN_SEARCH_MODE_ALL, then all non-null items in the index are considered candidate matches, whether they match any of the returned keys or not. (This mode is much slower than the other two choices, since it requires scanning essentially the entire index, but it may be necessary to implement corner cases correctly. An operator that needs this mode in most cases is probably not a good candidate for a GIN operator class.) The symbols to use for setting this mode are defined in access/gin.h.

pmatch is an output argument for use when partial match is supported. To use it, extractQuery must allocate an array of *nkeys booleans and store its address at *pmatch. Each element of the array should be set to TRUE if the corresponding key requires partial match, FALSE if not. If *pmatch is set to NULL then GIN assumes partial match is not required. The variable is initialized to NULL before call, so this argument can simply be ignored by operator classes that do not support partial match.

extra_data is an output argument that allows extractQuery to pass additional data to the consistent and comparePartial methods. To use it, extractQuery must allocate an array of *nkeys pointers and store its address at *extra_data, then store whatever it wants to into the individual pointers. The variable is initialized to NULL before call, so this argument can simply be ignored by operator classes that do not require extra data. If *extra_data is set, the whole array is passed to the consistent method, and the appropriate element to the comparePartial method.

An operator class must also provide a function to check if an indexed item matches the query. It comes in two flavors, a boolean consistent function, and a ternary triConsistent function. triConsistent covers the functionality of both, so providing triConsistent alone is sufficient. However, if the boolean variant is significantly cheaper to calculate, it can be advantageous to provide both. If only the boolean variant is provided, some optimizations that depend on refuting index items before fetching all the keys are disabled.
bool consistent(bool check[], StrategyNumber n, Datum query, int32 nkeys, Pointer extra_data[], bool *recheck, Datum queryKeys[], bool nullFlags[])

Returns TRUE if an indexed item satisfies the query operator with strategy number n (or might satisfy it, if the recheck indication is returned). This function does not have direct access to the indexed item’s value, since GIN does not store items explicitly. Rather, what is available is knowledge about which key values extracted from the query appear in a given indexed item. The check array has length nkeys, which is the same as the number of keys previously returned by extractQuery for this query datum. Each element of the check array is TRUE if the indexed item contains the corresponding query key, i.e., if (check[i] == TRUE) the i-th key of the extractQuery result array is present in the indexed item. The original query datum is passed in case the consistent method needs to consult it, and so are the queryKeys[] and nullFlags[] arrays previously returned by extractQuery. extra_data is the extra-data array returned by extractQuery, or NULL if none.

When extractQuery returns a null key in queryKeys[], the corresponding check[] element is TRUE if the indexed item contains a null key; that is, the semantics of check[] are like IS NOT DISTINCT FROM. The consistent function can examine the corresponding nullFlags[] element if it needs to tell the difference between a regular value match and a null match.

On success, *recheck should be set to TRUE if the heap tuple needs to be rechecked against the query operator, or FALSE if the index test is exact. That is, a FALSE return value guarantees that the heap tuple does not match the query; a TRUE return value with *recheck set to FALSE guarantees that the heap tuple does match the query; and a TRUE return value with *recheck set to TRUE means that the heap tuple might match the query, so it needs to be fetched and rechecked by evaluating the query operator directly against the originally indexed item.

GinTernaryValue triConsistent(GinTernaryValue check[], StrategyNumber n, Datum query, int32 nkeys, Pointer extra_data[], Datum queryKeys[], bool nullFlags[])

triConsistent is similar to consistent, but instead of booleans in the check vector, there are three possible values for each key: GIN_TRUE, GIN_FALSE and GIN_MAYBE. GIN_FALSE and GIN_TRUE have the same meaning as regular boolean values, while GIN_MAYBE means that the presence of that key is not known. When GIN_MAYBE values are present, the function should only return GIN_TRUE if the item certainly matches whether or not the index item contains the corresponding query keys. Likewise, the function must return GIN_FALSE only if the item certainly does not match, whether or not it contains the GIN_MAYBE keys. If the result depends on the GIN_MAYBE entries, i.e., the match cannot be confirmed or refuted based on the known query keys, the function must return GIN_MAYBE.

When there are no GIN_MAYBE values in the check vector, a GIN_MAYBE return value is the equivalent of setting the recheck flag in the boolean consistent function.

Optionally, an operator class for GIN can supply the following method:

int comparePartial(Datum partial_key, Datum key, StrategyNumber n, Pointer extra_data)

Compare a partial-match query key to an index key. Returns an integer whose sign indicates the result: less than zero means the index key does not match the query, but the index scan should continue; zero means that the index key does match the query; greater than zero indicates that the index scan should stop because no more matches are possible. The strategy number n of the operator that generated the partial match query is provided, in case its semantics are needed to determine when to end the scan. Also, extra_data is the corresponding element of the extra-data array made by extractQuery, or NULL if none. Null keys are never passed to this function.
To support “partial match” queries, an operator class must provide the `comparePartial` method, and its `extractQuery` method must set the `pmatch` parameter when a partial-match query is encountered. See Section 63.4.2 for details.

The actual data types of the various `Datum` values mentioned above vary depending on the operator class. The item values passed to `extractValue` are always of the operator class’s input type, and all key values must be of the class’s `STORAGE` type. The type of the `query` argument passed to `extractQuery`, `consistent` and `triConsistent` is whatever is the right-hand input type of the class member operator identified by the strategy number. This need not be the same as the indexed type, so long as key values of the correct type can be extracted from it. However, it is recommended that the SQL declarations of these three support functions use the opclass’s indexed data type for the `query` argument, even though the actual type might be something else depending on the operator.

63.4. Implementation

Internally, a GIN index contains a B-tree index constructed over keys, where each key is an element of one or more indexed items (a member of an array, for example) and where each tuple in a leaf page contains either a pointer to a B-tree of heap pointers (a “posting tree”), or a simple list of heap pointers (a “posting list”) when the list is small enough to fit into a single index tuple along with the key value.

As of PostgreSQL 9.1, null key values can be included in the index. Also, placeholder nulls are included in the index for indexed items that are null or contain no keys according to `extractValue`. This allows searches that should find empty items to do so.

Multicolumn GIN indexes are implemented by building a single B-tree over composite values (column number, key value). The key values for different columns can be of different types.

63.4.1. GIN Fast Update Technique

Updating a GIN index tends to be slow because of the intrinsic nature of inverted indexes: inserting or updating one heap row can cause many inserts into the index (one for each key extracted from the indexed item). As of PostgreSQL 8.4, GIN is capable of postponing much of this work by inserting new tuples into a temporary, unsorted list of pending entries. When the table is vacuumed or autoanalyzed, or when `gin_clean_pending_list` function is called, or if the pending list becomes larger than `gin_pending_list_limit`, the entries are moved to the main GIN data structure using the same bulk insert techniques used during initial index creation. This greatly improves GIN index update speed, even counting the additional vacuum overhead. Moreover the overhead work can be done by a background process instead of in foreground query processing.

The main disadvantage of this approach is that searches must scan the list of pending entries in addition to searching the regular index, and so a large list of pending entries will slow searches significantly. Another disadvantage is that, while most updates are fast, an update that causes the pending list to become “too large” will incur an immediate cleanup cycle and thus be much slower than other updates. Proper use of autovacuum can minimize both of these problems.

If consistent response time is more important than update speed, use of pending entries can be disabled by turning off the `fastupdate` storage parameter for a GIN index. See CREATE INDEX for details.
63.4.2. Partial Match Algorithm

GIN can support “partial match” queries, in which the query does not determine an exact match for one or more keys, but the possible matches fall within a reasonably narrow range of key values (within the key sorting order determined by the compare support method). The extractQuery method, instead of returning a key value to be matched exactly, returns a key value that is the lower bound of the range to be searched, and sets the pmatch flag true. The key range is then scanned using the comparePartial method. comparePartial must return zero for a matching index key, less than zero for a non-match that is still within the range to be searched, or greater than zero if the index key is past the range that could match.

63.5. GIN Tips and Tricks

Create vs. insert

Insertion into a GIN index can be slow due to the likelihood of many keys being inserted for each item. So, for bulk insertions into a table it is advisable to drop the GIN index and recreate it after finishing bulk insertion.

As of PostgreSQL 8.4, this advice is less necessary since delayed indexing is used (see Section 63.4.1 for details). But for very large updates it may still be best to drop and recreate the index.

maintenance_work_mem

Build time for a GIN index is very sensitive to the maintenance_work_mem setting; it doesn’t pay to skimp on work memory during index creation.

gin_pending_list_limit

During a series of insertions into an existing GIN index that has fastupdate enabled, the system will clean up the pending-entry list whenever the list grows larger than gin_pending_list_limit. To avoid fluctuations in observed response time, it’s desirable to have pending-list cleanup occur in the background (i.e., via autovacuum). Foreground cleanup operations can be avoided by increasing gin_pending_list_limit or making autovacuum more aggressive. However, enlarging the threshold of the cleanup operation means that if a foreground cleanup does occur, it will take even longer.

gin_pending_list_limit can be overridden for individual GIN indexes by changing storage parameters, and which allows each GIN index to have its own cleanup threshold. For example, it’s possible to increase the threshold only for the GIN index which can be updated heavily, and decrease it otherwise.

gin_fuzzy_search_limit

The primary goal of developing GIN indexes was to create support for highly scalable full-text search in PostgreSQL, and there are often situations when a full-text search returns a very large set of results. Moreover, this often happens when the query contains very frequent words, so that the large result set is not even useful. Since reading many tuples from the disk and sorting them could take a lot of time, this is unacceptable for production. (Note that the index search itself is very fast.)

To facilitate controlled execution of such queries, GIN has a configurable soft upper limit on the number of rows returned: the gin_fuzzy_search_limit configuration parameter. It is set to 0 (meaning no limit) by default. If a non-zero limit is set, then the returned set is a subset of the whole result set, chosen at random.
“Soft” means that the actual number of returned results could differ somewhat from the specified limit, depending on the query and the quality of the system’s random number generator. From experience, values in the thousands (e.g., 5000 — 20000) work well.

63.6. Limitations

GIN assumes that indexable operators are strict. This means that \texttt{extractValue} will not be called at all on a null item value (instead, a placeholder index entry is created automatically), and \texttt{extractQuery} will not be called on a null query value either (instead, the query is presumed to be unsatisfiable). Note however that null key values contained within a non-null composite item or query value are supported.

63.7. Examples

The PostgreSQL source distribution includes GIN operator classes for \texttt{tsvector} and for one-dimensional arrays of all internal types. Prefix searching in \texttt{tsvector} is implemented using the GIN partial match feature. The following \texttt{contrib} modules also contain GIN operator classes:

\begin{itemize}
  \item \texttt{btree_gin} \hfill \texttt{btree_gin}
    \begin{itemize}
      \item B-tree equivalent functionality for several data types
    \end{itemize}
  \item \texttt{hstore} \hfill \texttt{hstore}
    \begin{itemize}
      \item Module for storing (key, value) pairs
    \end{itemize}
  \item \texttt{intarray} \hfill \texttt{intarray}
    \begin{itemize}
      \item Enhanced support for \texttt{int[]}\n    \end{itemize}
  \item \texttt{pg_trgm} \hfill \texttt{pg_trgm}
    \begin{itemize}
      \item Text similarity using trigram matching
    \end{itemize}
\end{itemize}
Chapter 64. BRIN Indexes

64.1. Introduction

BRIN stands for Block Range Index. BRIN is designed for handling very large tables in which certain columns have some natural correlation with their physical location within the table. A block range is a group of pages that are physically adjacent in the table; for each block range, some summary info is stored by the index. For example, a table storing a store’s sale orders might have a date column on which each order was placed, and most of the time the entries for earlier orders will appear earlier in the table as well; a table storing a ZIP code column might have all codes for a city grouped together naturally.

BRIN indexes can satisfy queries via regular bitmap index scans, and will return all tuples in all pages within each range if the summary info stored by the index is consistent with the query conditions. The query executor is in charge of rechecking these tuples and discarding those that do not match the query conditions — in other words, these indexes are lossy. Because a BRIN index is very small, scanning the index adds little overhead compared to a sequential scan, but may avoid scanning large parts of the table that are known not to contain matching tuples.

The specific data that a BRIN index will store, as well as the specific queries that the index will be able to satisfy, depend on the operator class selected for each column of the index. Data types having a linear sort order can have operator classes that store the minimum and maximum value within each block range, for instance; geometrical types might store the bounding box for all the objects in the block range.

The size of the block range is determined at index creation time by the `pages_per_range` storage parameter. The number of index entries will be equal to the size of the relation in pages divided by the selected value for `pages_per_range`. Therefore, the smaller the number, the larger the index becomes (because of the need to store more index entries), but at the same time the summary data stored can be more precise and more data blocks can be skipped during an index scan.

64.1.1. Index Maintenance

At the time of creation, all existing index pages are scanned and a summary index tuple is created for each range, including the possibly-incomplete range at the end. As new pages are filled with data, page ranges that are already summarized will cause the summary information to be updated with data from the new tuples. When a new page is created that does not fall within the last summarized range, that range does not automatically acquire a summary tuple; those tuples remain unsummarized until a summarization run is invoked later, creating initial summaries. This process can be invoked manually using the `brin_summarize_new_values(regclass)` function, or automatically when `VACUUM` processes the table.

64.2. Built-in Operator Classes

The core PostgreSQL distribution includes the BRIN operator classes shown in Table 64-1.
The **minmax** operator classes store the minimum and the maximum values appearing in the indexed column within the range. The **inclusion** operator classes store a value which includes the values in the indexed column within the range.

Table 64-1. Built-in BRIN Operator Classes

<table>
<thead>
<tr>
<th>Name</th>
<th>Indexed Data Type</th>
<th>Indexable Operators</th>
</tr>
</thead>
<tbody>
<tr>
<td>abstime_minmax_ops</td>
<td>abstime</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>int8_minmax_ops</td>
<td>bigint</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>bit_minmax_ops</td>
<td>bit</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>varbit_minmax_ops</td>
<td>bit varying</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>box_inclusion_ops</td>
<td>box</td>
<td>&lt;&lt; &amp; &lt; &amp; &gt; &amp; &gt; = @ &gt;</td>
</tr>
<tr>
<td>bytea_minmax_ops</td>
<td>bytea</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>bpchar_minmax_ops</td>
<td>&quot;char&quot;</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>char_minmax_ops</td>
<td>date</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>float8_minmax_ops</td>
<td>double precision</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>inet_minmax_ops</td>
<td>inet</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>network_inclusion_ops</td>
<td>inet</td>
<td>&amp; &amp; &gt;&gt; &lt;&lt;&lt; = &lt;&lt;&lt;</td>
</tr>
<tr>
<td>int4_minmax_ops</td>
<td>integer</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>interval_minmax_ops</td>
<td>interval</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>macaddr_minmax_ops</td>
<td>macaddr</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>name_minmax_ops</td>
<td>name</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>numeric_minmax_ops</td>
<td>numeric</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>pg_lsn_minmax_ops</td>
<td>pg_lsn</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>oid_minmax_ops</td>
<td>oid</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>range_inclusion_ops</td>
<td>any range type</td>
<td>&lt;&lt; &amp; &lt;&lt; &amp; &gt;&gt; = @ = @</td>
</tr>
<tr>
<td>float4_minmax_ops</td>
<td>real</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>reltime_minmax_ops</td>
<td>reltime</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>int2_minmax_ops</td>
<td>smallint</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>text_minmax_ops</td>
<td>text</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>tid_minmax_ops</td>
<td>tid</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>timestamp_minmax_ops</td>
<td>timestamp without time zone</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>timestamptz_minmax_ops</td>
<td>timestamp with time zone</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>time_minmax_ops</td>
<td>time without time zone</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>timetz_minmax_ops</td>
<td>time with time zone</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
<tr>
<td>uuid_minmax_ops</td>
<td>uuid</td>
<td>&lt;&lt; &lt;= &gt;= &gt;&gt;</td>
</tr>
</tbody>
</table>
64.3. Extensibility

The BRIN interface has a high level of abstraction, requiring the access method implementer only to implement the semantics of the data type being accessed. The BRIN layer itself takes care of concurrency, logging and searching the index structure.

All it takes to get a BRIN access method working is to implement a few user-defined methods, which define the behavior of summary values stored in the index and the way they interact with scan keys. In short, BRIN combines extensibility with generality, code reuse, and a clean interface.

There are four methods that an operator class for BRIN must provide:

\begin{verbatim}
BrinOpcInfo *opcInfo(Oid type_oid)

Returns internal information about the indexed columns’ summary data. The return value must point to a palloc’d BrinOpcInfo, which has this definition:

typedef struct BrinOpcInfo
{
  /* Number of columns stored in an index column of this opclass */
  uint16   oi_nstored;

  /* Opaque pointer for the opclass’ private use */
  void *oi_opaque;

  /* Type cache entries of the stored columns */
  TypeCacheEntry *oi_typcache[FLEXIBLE_ARRAY_MEMBER];
} BrinOpcInfo;

BrinOpcInfo.oi_opaque can be used by the operator class routines to pass information between support procedures during an index scan.

bool consistent(BrinDesc *bdesc, BrinValues *column, ScanKey key)

Returns whether the ScanKey is consistent with the given indexed values for a range. The attribute number to use is passed as part of the scan key.

bool addValue(BrinDesc *bdesc, BrinValues *column, Datum newval, bool isnull)

Given an index tuple and an indexed value, modifies the indicated attribute of the tuple so that it additionally represents the new value. If any modification was done to the tuple, true is returned.

bool unionTuples(BrinDesc *bdesc, BrinValues *a, BrinValues *b)

Consolidates two index tuples. Given two index tuples, modifies the indicated attribute of the first of them so that it represents both tuples. The second tuple is not modified.
\end{verbatim}

The core distribution includes support for two types of operator classes: minmax and inclusion. Operator class definitions using them are shipped for in-core data types as appropriate. Additional operator classes can be defined by the user for other data types using equivalent definitions, without having to write any source code; appropriate catalog entries being declared is enough. Note that assumptions about the semantics of operator strategies are embedded in the support procedures’ source code.

Operator classes that implement completely different semantics are also possible, provided implementations of the four main support procedures described above are written. Note that backwards compatibility across major releases is not guaranteed: for example, additional support procedures might be required in later releases.

To write an operator class for a data type that implements a totally ordered set, it is possible to use the minmax support procedures alongside the corresponding operators, as shown in Table 64-2. All operator class members (procedures and operators) are mandatory.
Table 64-2. Procedure and Support Numbers for Minmax Operator Classes

<table>
<thead>
<tr>
<th>Operator class member</th>
<th>Object</th>
</tr>
</thead>
<tbody>
<tr>
<td>Support Procedure 1</td>
<td>internal function brin_minmax_opcinfo()</td>
</tr>
<tr>
<td>Support Procedure 2</td>
<td>internal function brin_minmax_add_value()</td>
</tr>
<tr>
<td>Support Procedure 3</td>
<td>internal function brin_minmax_consistent()</td>
</tr>
<tr>
<td>Support Procedure 4</td>
<td>internal function brin_minmax_union()</td>
</tr>
<tr>
<td>Operator Strategy 1</td>
<td>operator less-than</td>
</tr>
<tr>
<td>Operator Strategy 2</td>
<td>operator less-than-or-equal-to</td>
</tr>
<tr>
<td>Operator Strategy 3</td>
<td>operator equal-to</td>
</tr>
<tr>
<td>Operator Strategy 4</td>
<td>operator greater-than-or-equal-to</td>
</tr>
<tr>
<td>Operator Strategy 5</td>
<td>operator greater-than</td>
</tr>
</tbody>
</table>

To write an operator class for a complex data type which has values included within another type, it’s possible to use the inclusion support procedures alongside the corresponding operators, as shown in Table 64-3. It requires only a single additional function, which can be written in any language. More functions can be defined for additional functionality. All operators are optional. Some operators require other operators, as shown as dependencies on the table.

Table 64-3. Procedure and Support Numbers for Inclusion Operator Classes

<table>
<thead>
<tr>
<th>Operator class member</th>
<th>Object</th>
<th>Dependency</th>
</tr>
</thead>
<tbody>
<tr>
<td>Support Procedure 1</td>
<td>internal function brin_inclusion_opcinfo()</td>
<td></td>
</tr>
<tr>
<td>Support Procedure 2</td>
<td>internal function brin_inclusion_add_value()</td>
<td></td>
</tr>
<tr>
<td>Support Procedure 3</td>
<td>internal function brin_inclusion_consistent()</td>
<td></td>
</tr>
<tr>
<td>Support Procedure 4</td>
<td>internal function brin_inclusion_union()</td>
<td></td>
</tr>
<tr>
<td>Support Procedure 11</td>
<td>function to merge two elements</td>
<td></td>
</tr>
<tr>
<td>Support Procedure 12</td>
<td>optional function to check whether two elements are mergeable</td>
<td>Operator Strategy 4</td>
</tr>
<tr>
<td>Support Procedure 13</td>
<td>optional function to check if an element is contained within another</td>
<td>Operator Strategy 5</td>
</tr>
<tr>
<td>Support Procedure 14</td>
<td>optional function to check whether an element is empty</td>
<td>Operator Strategy 4</td>
</tr>
<tr>
<td>Operator Strategy 1</td>
<td>operator left-of</td>
<td>Operator Strategy 4</td>
</tr>
<tr>
<td>Operator Strategy 3</td>
<td>operator overlaps</td>
<td></td>
</tr>
</tbody>
</table>
## Operator class member

<table>
<thead>
<tr>
<th>Operator Strategy</th>
<th>Object</th>
<th>Dependency</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td>operator does-not-extend-to-the-left-of</td>
<td>Operator Strategy 1</td>
</tr>
<tr>
<td>5</td>
<td>operator right-of</td>
<td>Operator Strategy 2</td>
</tr>
<tr>
<td>6, 18</td>
<td>operator same-as-or-equal-to</td>
<td>Operator Strategy 7</td>
</tr>
<tr>
<td>7, 13, 16, 24, 25</td>
<td>operator contains-or-equal-to</td>
<td>Operator Strategy 3</td>
</tr>
<tr>
<td>8, 14, 26, 27</td>
<td>operator is-contained-by-or-equal-to</td>
<td>Operator Strategy 11</td>
</tr>
<tr>
<td>9</td>
<td>operator does-not-extend-above</td>
<td>Operator Strategy 12</td>
</tr>
<tr>
<td>10</td>
<td>operator is-below</td>
<td>Operator Strategy 12</td>
</tr>
<tr>
<td>11</td>
<td>operator is-above</td>
<td>Operator Strategy 9</td>
</tr>
<tr>
<td>12</td>
<td>operator does-not-extend-below</td>
<td>Operator Strategy 10</td>
</tr>
<tr>
<td>20</td>
<td>operator less-than</td>
<td>Operator Strategy 5</td>
</tr>
<tr>
<td>21</td>
<td>operator less-than-or-equal-to</td>
<td>Operator Strategy 5</td>
</tr>
<tr>
<td>22</td>
<td>operator greater-than</td>
<td>Operator Strategy 1</td>
</tr>
<tr>
<td>23</td>
<td>operator greater-than-or-equal-to</td>
<td>Operator Strategy 1</td>
</tr>
</tbody>
</table>

Support procedure numbers 1-10 are reserved for the BRIN internal functions, so the SQL level functions start with number 11. Support function number 11 is the main function required to build the index. It should accept two arguments with the same data type as the operator class, and return the union of them. The inclusion operator class can store union values with different data types if it is defined with the `STORAGE` parameter. The return value of the union function should match the `STORAGE` data type.

Support procedure numbers 12 and 14 are provided to support irregularities of built-in data types. Procedure number 12 is used to support network addresses from different families which are not mergeable. Procedure number 14 is used to support empty ranges. Procedure number 13 is an optional but recommended one, which allows the new value to be checked before it is passed to the union function. As the BRIN framework can shortcut some operations when the union is not changed, using this function can improve index performance.

Both minmax and inclusion operator classes support cross-data-type operators, though with these the dependencies become more complicated. The minmax operator class requires a full set of operators to be defined with both arguments having the same data type. It allows additional data types to be supported by defining extra sets of operators. Inclusion operator class operator strategies are dependent on another operator strategy as shown in Table 64-3, or the same operator strategy as themselves. They require the dependency operator to be defined with the `STORAGE` data type as the left-hand-side argument and the other supported data type to be the right-hand-side argument of the supported operator. See `float4_minmax_ops` as an example of minmax, and `box_inclusion_ops` as an example of inclusion.
Chapter 65. Database Physical Storage

This chapter provides an overview of the physical storage format used by PostgreSQL databases.

65.1. Database File Layout

This section describes the storage format at the level of files and directories.

Traditionally, the configuration and data files used by a database cluster are stored together within the cluster’s data directory, commonly referred to as PGDATA (after the name of the environment variable that can be used to define it). A common location for PGDATA is /var/lib/pgsql/data. Multiple clusters, managed by different server instances, can exist on the same machine.

The PGDATA directory contains several subdirectories and control files, as shown in Table 65-1. In addition to these required items, the cluster configuration files postgresql.conf, pg_hba.conf, and pg_ident.conf are traditionally stored in PGDATA, although it is possible to place them elsewhere.

Table 65-1. Contents of PGDATA

<table>
<thead>
<tr>
<th>Item</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PG_VERSION</td>
<td>A file containing the major version number of PostgreSQL</td>
</tr>
<tr>
<td>base</td>
<td>Subdirectory containing per-database subdirectories</td>
</tr>
<tr>
<td>global</td>
<td>Subdirectory containing cluster-wide tables, such as pg_database</td>
</tr>
<tr>
<td>pg_commit_ts</td>
<td>Subdirectory containing transaction commit timestamp data</td>
</tr>
<tr>
<td>pg_clog</td>
<td>Subdirectory containing transaction commit status data</td>
</tr>
<tr>
<td>pg_dynshmem</td>
<td>Subdirectory containing files used by the dynamic shared memory subsystem</td>
</tr>
<tr>
<td>pg_logical</td>
<td>Subdirectory containing status data for logical decoding</td>
</tr>
<tr>
<td>pg_multixact</td>
<td>Subdirectory containing multitransaction status data (used for shared row locks)</td>
</tr>
<tr>
<td>pg_notify</td>
<td>Subdirectory containing LISTEN/NOTIFY status data</td>
</tr>
<tr>
<td>pg_replslot</td>
<td>Subdirectory containing replication slot data</td>
</tr>
<tr>
<td>pg_serial</td>
<td>Subdirectory containing information about committed serializable transactions</td>
</tr>
<tr>
<td>pg_snapshots</td>
<td>Subdirectory containing exported snapshots</td>
</tr>
<tr>
<td>pg_stat</td>
<td>Subdirectory containing permanent files for the statistics subsystem</td>
</tr>
<tr>
<td>Item</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td><code>pg_stat_tmp</code></td>
<td>Subdirectory containing temporary files for the statistics subsystem</td>
</tr>
<tr>
<td><code>pg_subtrans</code></td>
<td>Subdirectory containing subtransaction status data</td>
</tr>
<tr>
<td><code>pg_tblspc</code></td>
<td>Subdirectory containing symbolic links to tablespaces</td>
</tr>
<tr>
<td><code>pg_twophase</code></td>
<td>Subdirectory containing state files for prepared transactions</td>
</tr>
<tr>
<td><code>pg_xlog</code></td>
<td>Subdirectory containing WAL (Write Ahead Log) files</td>
</tr>
<tr>
<td><code>postgresql.auto.conf</code></td>
<td>A file used for storing configuration parameters that are set by \textit{ALTER SYSTEM}</td>
</tr>
<tr>
<td><code>postmaster.opts</code></td>
<td>A file recording the command-line options the server was last started with</td>
</tr>
<tr>
<td><code>postmaster.pid</code></td>
<td>A lock file recording the current postmaster process ID (PID), cluster data directory path, postmaster start timestamp, port number, Unix-domain socket directory path (empty on Windows), first valid listen_address (IP address or *), or empty if not listening on TCP, and shared memory segment ID (this file is not present after server shutdown)</td>
</tr>
</tbody>
</table>

For each database in the cluster there is a subdirectory within \texttt{PGDATA/base}, named after the database’s OID in \texttt{pg_database}. This subdirectory is the default location for the database’s files; in particular, its system catalogs are stored there.

Each table and index is stored in a separate file. For ordinary relations, these files are named after the table or index’s \texttt{filenode} number, which can be found in \texttt{pg_class}.\texttt{relfilenode}. But for temporary relations, the file name is of the form \texttt{tBBB_Find}, where \texttt{BBB} is the backend ID of the backend which created the file, and \texttt{FFF} is the filenode number. In either case, in addition to the main file (a/k/a main fork), each table and index has a free space map (see Section 65.3), which stores information about free space available in the relation. The free space map is stored in a file named with the filenode number plus the suffix \texttt{_fsm}. Tables also have a visibility map, stored in a fork with the suffix \texttt{_vm}, to track which pages are known to have no dead tuples. The visibility map is described further in Section 65.4. Unlogged tables and indexes have a third fork, known as the initialization fork, which is stored in a fork with the suffix \texttt{_init} (see Section 65.5).

\begin{caution}

Note that while a table’s filenode often matches its OID, this is \textit{not} necessarily the case; some operations, like \texttt{TRUNCATE}, \texttt{REINDEX}, \texttt{CLUSTER} and some forms of \texttt{ALTER TABLE}, can change the filenode while preserving the OID. Avoid assuming that filenode and table OID are the same. Also, for certain system catalogs including \texttt{pg_class} itself, \texttt{pg_class.relfilenode} contains zero. The actual filenode number of these catalogs is stored in a lower-level data structure, and can be obtained using the \texttt{pg_relation_filenode()} function.

\end{caution}

When a table or index exceeds 1 GB, it is divided into gigabyte-sized \textit{segments}. The first segment’s file name is the same as the filenode; subsequent segments are named filenode.1, filenode.2,
etc. This arrangement avoids problems on platforms that have file size limitations. (Actually, 1 GB is just the default segment size. The segment size can be adjusted using the configuration option `--with-segsize` when building PostgreSQL.) In principle, free space map and visibility map forks could require multiple segments as well, though this is unlikely to happen in practice.

A table that has columns with potentially large entries will have an associated TOAST table, which is used for out-of-line storage of field values that are too large to keep in the table rows proper. `pg_class.reltoastrelid` links from a table to its TOAST table, if any. See Section 65.2 for more information.

The contents of tables and indexes are discussed further in Section 65.6.

Tablespaces make the scenario more complicated. Each user-defined tablespace has a symbolic link inside the `PGDATA/pg_tblspc` directory, which points to the physical tablespace directory (i.e., the location specified in the tablespace’s `CREATE TABLESPACE` command). This symbolic link is named after the tablespace’s OID. Inside the physical tablespace directory there is a subdirectory with a name that depends on the PostgreSQL server version, such as `PG_9.0_20100805`. (The reason for using this subdirectory is so that successive versions of the database can use the same `CREATE TABLESPACE` location value without conflicts.) Within the version-specific subdirectory, there is a subdirectory for each database that has elements in the tablespace, named after the database’s OID. Tables and indexes are stored within that directory, using the filenode naming scheme. The `pg_default` tablespace is not accessed through `pg_tblspc`, but corresponds to `PGDATA/base`. Similarly, the `pg_global` tablespace is not accessed through `pg_tblspc`, but corresponds to `PGDATA/global`.

The `pg_relation_filepath()` function shows the entire path (relative to `PGDATA`) of any relation. It is often useful as a substitute for remembering many of the above rules. But keep in mind that this function just gives the name of the first segment of the main fork of the relation — you may need to append a segment number and/or `_fsm`, `_vm`, or `_init` to find all the files associated with the relation.

Temporary files (for operations such as sorting more data than can fit in memory) are created within `PGDATA/base/pgsql_tmp`, or within a `pgsql_tmp` subdirectory of a tablespace directory if a tablespace other than `pg_default` is specified for them. The name of a temporary file has the form `pgsql_tmpPPP_NNN`, where `PPP` is the PID of the owning backend and `NNN` distinguishes different temporary files of that backend.

## 65.2. TOAST

This section provides an overview of TOAST (The Oversized-Attribute Storage Technique).

PostgreSQL uses a fixed page size (commonly 8 kB), and does not allow tuples to span multiple pages. Therefore, it is not possible to store very large field values directly. To overcome this limitation, large field values are compressed and/or broken up into multiple physical rows. This happens transparently to the user, with only small impact on most of the backend code. The technique is affectionately known as TOAST (or “the best thing since sliced bread”). The TOAST infrastructure is also used to improve handling of large data values in-memory.

Only certain data types support TOAST — there is no need to impose the overhead on data types that cannot produce large field values. To support TOAST, a data type must have a variable-length (`varlena`) representation, in which, ordinarily, the first four-byte word of any stored value contains the total length of the value in bytes (including itself). TOAST does not constrain the rest of the data type’s representation. The special representations collectively called `TOASTed values` work by modifying or reinterpreting this initial length word. Therefore, the C-level functions supporting a TOAST-able data type must be careful about how they handle potentially TOASTed input values: an input might not actually consist of a four-byte length word and contents until after it’s been `detoasted`.

2183
(This is normally done by invoking `PG_DETOAST_DATUM` before doing anything with an input value, but in some cases more efficient approaches are possible. See Section 36.11.1 for more detail.)

TOAST usurps two bits of the varlena length word (the high-order bits on big-endian machines, the low-order bits on little-endian machines), thereby limiting the logical size of any value of a TOAST-able data type to 1 GB ($2^{30} - 1$ bytes). When both bits are zero, the value is an ordinary un-TOASTed value of the data type, and the remaining bits of the length word give the total datum size (including length word) in bytes. When the highest-order or lowest-order bit is set, the value has only a single-byte header instead of the normal four-byte header, and the remaining bits of that byte give the total datum size (including length byte) in bytes. This alternative supports space-efficient storage of values shorter than 127 bytes, while still allowing the data type to grow to 1 GB at need. Values with single-byte headers aren’t aligned on any particular boundary, whereas values with four-byte headers are aligned on at least a four-byte boundary; this omission of alignment padding provides additional space savings that is significant compared to short values. As a special case, if the remaining bits of a single-byte header are all zero (which would be impossible for a self-inclusive length), the value is a pointer to out-of-line data, with several possible alternatives as described below. The type and size of such a TOAST pointer are determined by a code stored in the second byte of the datum. Lastly, when the highest-order or lowest-order bit is clear but the adjacent bit is set, the content of the datum has been compressed and must be decompressed before use. In this case the remaining bits of the four-byte length word give the total size of the compressed datum, not the original data. Note that compression is also possible for out-of-line data but the varlena header does not tell whether it has occurred — the content of the TOAST pointer tells that, instead.

As mentioned, there are multiple types of TOAST pointer datums. The oldest and most common type is a pointer to out-of-line data stored in a TOAST table that is separate from, but associated with, the table containing the TOAST pointer datum itself. These on-disk pointer datums are created by the TOAST management code (in `access/heap/tuptoaster.c`) when a tuple to be stored on disk is too large to be stored as-is. Further details appear in Section 65.2.1. Alternatively, a TOAST pointer datum can contain a pointer to out-of-line data that appears elsewhere in memory. Such datums are necessarily short-lived, and will never appear on-disk, but they are very useful for avoiding copying and redundant processing of large data values. Further details appear in Section 65.2.2.

The compression technique used for either in-line or out-of-line compressed data is a fairly simple and very fast member of the LZ family of compression techniques. See `src/common/pg_lzcompress.c` for the details.

### 65.2.1. Out-of-line, on-disk TOAST storage

If any of the columns of a table are TOAST-able, the table will have an associated TOAST table, whose OID is stored in the table’s `pg_class.reltoastrelid` entry. On-disk TOASTed values are kept in the TOAST table, as described in more detail below.

Out-of-line values are divided (after compression if used) into chunks of at most `TOAST_MAX_CHUNK_SIZE` bytes (by default this value is chosen so that four chunk rows will fit on a page, making it about 2000 bytes). Each chunk is stored as a separate row in the TOAST table belonging to the owning table. Every TOAST table has the columns `chunk_id` (an OID identifying the particular TOASTed value), `chunk_seq` (a sequence number for the chunk within its value), and `chunk_data` (the actual data of the chunk). A unique index on `chunk_id` and `chunk_seq` provides fast retrieval of the values. A pointer datum representing an out-of-line on-disk TOASTed value therefore needs to store the OID of the TOAST table in which to look and the OID of the specific value (its `chunk_id`). For convenience, pointer datums also store the logical datum size (original uncompressed data length) and physical stored size (different if compression was applied). Allowing for the varlena header bytes, the total size of an on-disk TOAST pointer datum is therefore 18 bytes.
regardless of the actual size of the represented value.

The TOAST management code is triggered only when a row value to be stored in a table is wider than `TOAST_TUPLE_THRESHOLD` bytes (normally 2 kB). The TOAST code will compress and/or move field values out-of-line until the row value is shorter than `TOAST_TUPLE_TARGET` bytes (also normally 2 kB) or no more gains can be had. During an UPDATE operation, values of unchanged fields are normally preserved as-is; so an UPDATE of a row with out-of-line values incurs no TOAST costs if none of the out-of-line values change.

The TOAST management code recognizes four different strategies for storing TOAST-able columns on disk:

- **PLAIN** prevents either compression or out-of-line storage; furthermore it disables use of single-byte headers for varlena types. This is the only possible strategy for columns of non-TOAST-able data types.
- **EXTENDED** allows both compression and out-of-line storage. This is the default for most TOAST-able data types. Compression will be attempted first, then out-of-line storage if the row is still too big.
- **EXTERNAL** allows out-of-line storage but not compression. Use of **EXTERNAL** will make substring operations on wide `text` and `bytea` columns faster (at the penalty of increased storage space) because these operations are optimized to fetch only the required parts of the out-of-line value when it is not compressed.
- **MAIN** allows compression but not out-of-line storage. (Actually, out-of-line storage will still be performed for such columns, but only as a last resort when there is no other way to make the row small enough to fit on a page.)

Each TOAST-able data type specifies a default strategy for columns of that data type, but the strategy for a given table column can be altered with `ALTER TABLE SET STORAGE`.

This scheme has a number of advantages compared to a more straightforward approach such as allowing row values to span pages. Assuming that queries are usually qualified by comparisons against relatively small key values, most of the work of the executor will be done using the main row entry. The big values of TOASTed attributes will only be pulled out (if selected at all) at the time the result set is sent to the client. Thus, the main table is much smaller and more of its rows fit in the shared buffer cache than would be the case without any out-of-line storage. Sort sets shrink also, and sorts will more often be done entirely in memory. A little test showed that a table containing typical HTML pages and their URLs was stored in about half of the raw data size including the TOAST table, and that the main table contained only about 10% of the entire data (the URLs and some small HTML pages). There was no run time difference compared to an un-TOASTed comparison table, in which all the HTML pages were cut down to 7 kB to fit.

### 65.2.2. Out-of-line, in-memory TOAST storage

TOAST pointers can point to data that is not on disk, but is elsewhere in the memory of the current server process. Such pointers obviously cannot be long-lived, but they are nonetheless useful. There are currently two sub-cases: pointers to **indirect** data and pointers to **expanded** data.

Indirect TOAST pointers simply point at a non-indirect varlena value stored somewhere in memory. This case was originally created merely as a proof of concept, but it is currently used during logical decoding to avoid possibly having to create physical tuples exceeding 1 GB (as pulling all out-of-line field values into the tuple might do). The case is of limited use since the creator of the pointer datum
is entirely responsible that the referenced data survives for as long as the pointer could exist, and there
is no infrastructure to help with this.

Expanded TOAST pointers are useful for complex data types whose on-disk representation is not es-
pecially suited for computational purposes. As an example, the standard varlena representation of a
PostgreSQL array includes dimensionality information, a nulls bitmap if there are any null elements,
then the values of all the elements in order. When the element type itself is variable-length, the only
way to find the \( N \)’th element is to scan through all the preceding elements. This representation is
appropriate for on-disk storage because of its compactness, but for computations with the array it’s
much nicer to have an “expanded” or “deconstructed” representation in which all the element start-
ing locations have been identified. The TOAST pointer mechanism supports this need by allowing a
pass-by-reference Datum to point to either a standard varlena value (the on-disk representation) or
a TOAST pointer that points to an expanded representation somewhere in memory. The details of
this expanded representation are up to the data type, though it must have a standard header and meet
the other API requirements given in src/include/utils/expandeddatum.h. C-level functions
working with the data type can choose to handle either representation. Functions that do not know
about the expanded representation, but simply apply \texttt{PG\_DETOAST\_DATUM} to their inputs, will auto-
matically receive the traditional varlena representation; so support for an expanded representation can
be introduced incrementally, one function at a time.

TOAST pointers to expanded values are further broken down into \texttt{read-write} and \texttt{read-only}
pointers. The pointed-to representation is the same either way, but a function that receives a read-write pointer
is allowed to modify the referenced value in-place, whereas one that receives a read-only pointer must
not; it must first create a copy if it wants to make a modified version of the value. This distinction
and some associated conventions make it possible to avoid unnecessary copying of expanded values
during query execution.

For all types of in-memory TOAST pointer, the TOAST management code ensures that no such
pointer datum can accidentally get stored on disk. In-memory TOAST pointers are automatically
expanded to normal in-line varlena values before storage — and then possibly converted to on-disk
TOAST pointers, if the containing tuple would otherwise be too big.

\section*{65.3. Free Space Map}

Each heap and index relation, except for hash indexes, has a Free Space Map (FSM) to keep track of
available space in the relation. It’s stored alongside the main relation data in a separate relation fork,
named after the filenode number of the relation, plus a \texttt{.fsm} suffix. For example, if the filenode of
a relation is 12345, the FSM is stored in a file called \texttt{12345.fsm}, in the same directory as the main
relation file.

The Free Space Map is organized as a tree of FSM pages. The bottom level FSM pages store the free
space available on each heap (or index) page, using one byte to represent each such page. The upper
levels aggregate information from the lower levels.

Within each FSM page is a binary tree, stored in an array with one byte per node. Each leaf node
represents a heap page, or a lower level FSM page. In each non-leaf node, the higher of its children’s
values is stored. The maximum value in the leaf nodes is therefore stored at the root.

See src/backend/storage/freespace/README for more details on how the FSM is structured,
and how it’s updated and searched. The \texttt{pg\_freespacemap} module can be used to examine the infor-
mation stored in free space maps.
65.4. Visibility Map

Each heap relation has a Visibility Map (VM) to keep track of which pages contain only tuples that are known to be visible to all active transactions; it also keeps track of which pages contain only frozen tuples. It’s stored alongside the main relation data in a separate relation fork, named after the filenode number of the relation, plus a _vm suffix. For example, if the filenode of a relation is 12345, the VM is stored in a file called 12345_vm, in the same directory as the main relation file. Note that indexes do not have VMs.

The visibility map stores two bits per heap page. The first bit, if set, indicates that the page is all-visible, or in other words that the page does not contain any tuples that need to be vacuumed. This information can also be used by index-only scans to answer queries using only the index tuple. The second bit, if set, means that all tuples on the page have been frozen. That means that even an anti-wraparound vacuum need not revisit the page.

The map is conservative in the sense that we make sure that whenever a bit is set, we know the condition is true, but if a bit is not set, it might or might not be true. Visibility map bits are only set by vacuum, but are cleared by any data-modifying operations on a page.

The pg_visibility module can be used to examine the information stored in the visibility map.

65.5. The Initialization Fork

Each unlogged table, and each index on an unlogged table, has an initialization fork. The initialization fork is an empty table or index of the appropriate type. When an unlogged table must be reset to empty due to a crash, the initialization fork is copied over the main fork, and any other forks are erased (they will be recreated automatically as needed).

65.6. Database Page Layout

This section provides an overview of the page format used within PostgreSQL tables and indexes. Sequences and TOAST tables are formatted just like a regular table.

In the following explanation, a byte is assumed to contain 8 bits. In addition, the term item refers to an individual data value that is stored on a page. In a table, an item is a row; in an index, an item is an index entry.

Every table and index is stored as an array of pages of a fixed size (usually 8 kB, although a different page size can be selected when compiling the server). In a table, all the pages are logically equivalent, so a particular item (row) can be stored in any page. In indexes, the first page is generally reserved as a metapage holding control information, and there can be different types of pages within the index, depending on the index access method.

Table 65-2 shows the overall layout of a page. There are five parts to each page.

Table 65-2. Overall Page Layout

<table>
<thead>
<tr>
<th>Item</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PageHeaderData</td>
<td>24 bytes long. Contains general information about the page, including free space pointers.</td>
</tr>
</tbody>
</table>

1. Actually, index access methods need not use this page format. All the existing index methods do use this basic format, but the data kept on index metapages usually doesn’t follow the item layout rules.
## Chapter 65. Database Physical Storage

<table>
<thead>
<tr>
<th>Item</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ItemIdData</td>
<td>Array of item identifiers pointing to the actual items. Each entry is an (offset,length) pair. 4 bytes per item.</td>
</tr>
<tr>
<td>Free space</td>
<td>The unallocated space. New item identifiers are allocated from the start of this area, new items from the end.</td>
</tr>
<tr>
<td>Items</td>
<td>The actual items themselves.</td>
</tr>
<tr>
<td>Special space</td>
<td>Index access method specific data. Different methods store different data. Empty in ordinary tables.</td>
</tr>
</tbody>
</table>

The first 24 bytes of each page consists of a page header (**PageHeaderData**). Its format is detailed in Table 65-3. The first field tracks the most recent WAL entry related to this page. The second field contains the page checksum if data checksums are enabled. Next is a 2-byte field containing flag bits. This is followed by three 2-byte integer fields (**pd_lower**, **pd_upper**, and **pd_special**). These contain byte offsets from the page start to the start of unallocated space, to the end of unallocated space, and to the start of the special space. The next 2 bytes of the page header, **pd_pagesize_version**, store both the page size and a version indicator. Beginning with PostgreSQL 8.3 the version number is 4; PostgreSQL 8.1 and 8.2 used version number 3; PostgreSQL 8.0 used version number 2; PostgreSQL 7.3 and 7.4 used version number 1; prior releases used version number 0. (The basic page layout and header format has not changed in most of these versions, but the layout of heap row headers has.) The page size is basically only present as a cross-check; there is no support for having more than one page size in an installation. The last field is a hint that shows whether pruning the page is likely to be profitable: it tracks the oldest un-pruned XMAX on the page.

**Table 65-3. PageHeaderData Layout**

<table>
<thead>
<tr>
<th>Field</th>
<th>Type</th>
<th>Length</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>pd_lsn</td>
<td>PageXLogRecPtr</td>
<td>8 bytes</td>
<td>LSN: next byte after last byte of xlog record for last change to this page</td>
</tr>
<tr>
<td>pd_checksum</td>
<td>uint16</td>
<td>2 bytes</td>
<td>Page checksum</td>
</tr>
<tr>
<td>pd_flags</td>
<td>uint16</td>
<td>2 bytes</td>
<td>Flag bits</td>
</tr>
<tr>
<td>pd_lower</td>
<td>LocationIndex</td>
<td>2 bytes</td>
<td>Offset to start of free space</td>
</tr>
<tr>
<td>pd_upper</td>
<td>LocationIndex</td>
<td>2 bytes</td>
<td>Offset to end of free space</td>
</tr>
<tr>
<td>pd_special</td>
<td>LocationIndex</td>
<td>2 bytes</td>
<td>Offset to start of special space</td>
</tr>
<tr>
<td>pd_pagesize_version</td>
<td>uint16</td>
<td>2 bytes</td>
<td>Page size and layout version number information</td>
</tr>
<tr>
<td>pd_prune_xid</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>Oldest unpruned XMAX on page, or zero if none</td>
</tr>
</tbody>
</table>

All the details can be found in **src/include/storage/bufpage.h**.
Following the page header are item identifiers (ItemIdData), each requiring four bytes. An item identifier contains a byte-offset to the start of an item, its length in bytes, and a few attribute bits which affect its interpretation. New item identifiers are allocated as needed from the beginning of the unallocated space. The number of item identifiers present can be determined by looking at \texttt{pd_lower}, which is increased to allocate a new identifier. Because an item identifier is never moved until it is freed, its index can be used on a long-term basis to reference an item, even when the item itself is moved around on the page to compact free space. In fact, every pointer to an item (ItemIdPointer, also known as CTID) created by PostgreSQL consists of a page number and the index of an item identifier.

The items themselves are stored in space allocated backwards from the end of unallocated space. The exact structure varies depending on what the table is to contain. Tables and sequences both use a structure named \texttt{HeapTupleHeaderData}, described below.

The final section is the “special section” which can contain anything the access method wishes to store. For example, b-tree indexes store links to the page’s left and right siblings, as well as some other data relevant to the index structure. Ordinary tables do not use a special section at all (indicated by setting \texttt{pd_special} to equal the page size).

All table rows are structured in the same way. There is a fixed-size header (occupying 23 bytes on most machines), followed by an optional null bitmap, an optional object ID field, and the user data. The header is detailed in Table 65-4. The actual user data (columns of the row) begins at the offset indicated by \texttt{t_hoff}, which must always be a multiple of the MAXALIGN distance for the platform. The null bitmap is only present if the \texttt{HEAP_HASNULL} bit is set in \texttt{t_infomask}. If it is present it begins just after the fixed header and occupies enough bytes to have one bit per data column (that is, \texttt{t_natts} bits altogether). In this list of bits, a 1 bit indicates not-null, a 0 bit is a null. When the bitmap is not present, all columns are assumed not-null. The object ID is only present if the \texttt{HEAP_HASOID} bit is set in \texttt{t_infomask}. If present, it appears just before the \texttt{t_hoff} boundary. Any padding needed to make \texttt{t_hoff} a MAXALIGN multiple will appear between the null bitmap and the object ID. (This in turn ensures that the object ID is suitably aligned.)

### Table 65-4. HeapTupleHeaderData Layout

<table>
<thead>
<tr>
<th>Field</th>
<th>Type</th>
<th>Length</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>\texttt{t_xmin}</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>insert XID stamp</td>
</tr>
<tr>
<td>\texttt{t_xmax}</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>delete XID stamp</td>
</tr>
<tr>
<td>\texttt{t_cid}</td>
<td>CommandId</td>
<td>4 bytes</td>
<td>insert and/or delete CID stamp (overlays with \texttt{t_xvac})</td>
</tr>
<tr>
<td>\texttt{t_xvac}</td>
<td>TransactionId</td>
<td>4 bytes</td>
<td>XID for VACUUM operation moving a row version</td>
</tr>
<tr>
<td>\texttt{t_ctid}</td>
<td>ItemPointerData</td>
<td>6 bytes</td>
<td>current TID of this or newer row version</td>
</tr>
<tr>
<td>\texttt{t_infomask2}</td>
<td>uint16</td>
<td>2 bytes</td>
<td>number of attributes, plus various flag bits</td>
</tr>
<tr>
<td>\texttt{t_infomask}</td>
<td>uint16</td>
<td>2 bytes</td>
<td>various flag bits</td>
</tr>
<tr>
<td>\texttt{t_hoff}</td>
<td>uint8</td>
<td>1 byte</td>
<td>offset to user data</td>
</tr>
</tbody>
</table>

All the details can be found in \texttt{src/include/access/htup_details.h}.

Interpreting the actual data can only be done with information obtained from other tables, mostly \texttt{pg_attribute}. The key values needed to identify field locations are \texttt{attlen} and \texttt{attalign}. There
is no way to directly get a particular attribute, except when there are only fixed width fields and no null values. All this trickery is wrapped up in the functions \texttt{heapgetattr, fastgetattr} and \texttt{heap_getsysattr}.

To read the data you need to examine each attribute in turn. First check whether the field is NULL according to the null bitmap. If it is, go to the next. Then make sure you have the right alignment. If the field is a fixed width field, then all the bytes are simply placed. If it’s a variable length field (attlen = -1) then it’s a bit more complicated. All variable-length data types share the common header \texttt{struct varlena}, which includes the total length of the stored value and some flag bits. Depending on the flags, the data can be either inline or in a TOAST table; it might be compressed, too (see Section 65.2).
Chapter 66. BKI Backend Interface

Backend Interface (BKI) files are scripts in a special language that is understood by the PostgreSQL backend when running in the “bootstrap” mode. The bootstrap mode allows system catalogs to be created and filled from scratch, whereas ordinary SQL commands require the catalogs to exist already. BKI files can therefore be used to create the database system in the first place. (And they are probably not useful for anything else.)

initdb uses a BKI file to do part of its job when creating a new database cluster. The input file used by initdb is created as part of building and installing PostgreSQL by a program named genbki.pl, which reads some specially formatted C header files in the src/include/catalog/ directory of the source tree. The created BKI file is called postgres.bki and is normally installed in the share subdirectory of the installation tree.

Related information can be found in the documentation for initdb.

66.1. BKI File Format

This section describes how the PostgreSQL backend interprets BKI files. This description will be easier to understand if the postgres.bki file is at hand as an example.

BKI input consists of a sequence of commands. Commands are made up of a number of tokens, depending on the syntax of the command. Tokens are usually separated by whitespace, but need not be if there is no ambiguity. There is no special command separator; the next token that syntactically cannot belong to the preceding command starts a new one. (Usually you would put a new command on a new line, for clarity.) Tokens can be certain key words, special characters (parentheses, commas, etc.), numbers, or double-quoted strings. Everything is case sensitive.

Lines starting with # are ignored.

66.2. BKI Commands

```
create  tablename  tableoid  [bootstrap]  [shared_relation]  [without_oids]  
        [rowtype_oid  oid]  (name1  =  type1  [FORCE NOT NULL  |  FORCE NULL]  [,  name2  =  type2  
        [FORCE NOT NULL  |  FORCE NULL]  ],  ...])
```

Create a table named `tablename`, and having the OID `tableoid`, with the columns given in parentheses.

The following column types are supported directly by `bootstrap.c`: `bool`, `bytea`, `char` (1 byte), `name`, `int2`, `int4`, `regproc`, `regclass`, `regtype`, `text`, `oid`, `tid`, `xid`, `cid`, `int2vector`, `oidvector`, `_int4` (array), `_text` (array), `_oid` (array), `_char` (array), `_aclitem` (array). Although it is possible to create tables containing columns of other types, this cannot be done until after `pg_type` has been created and filled with appropriate entries. (That effectively means that only these column types can be used in bootstrapped tables, but non-bootstrap catalogs can contain any built-in type.)

When `bootstrap` is specified, the table will only be created on disk; nothing is entered into `pg_class`, `pg_attribute`, etc. for it. Thus the table will not be accessible by ordinary SQL
operations until such entries are made the hard way (with insert commands). This option is used for creating pg_class etc themselves.

The table is created as shared if shared_relation is specified. It will have OIDs unless without_oids is specified. The table’s row type OID (pg_type OID) can optionally be specified via the rowtype_oid clause; if not specified, an OID is automatically generated for it. (The rowtype_oid clause is useless if bootstrap is specified, but it can be provided anyway for documentation.)

**open tablename**

Open the table named tablename for insertion of data. Any currently open table is closed.

**close [tablename]**

Close the open table. The name of the table can be given as a cross-check, but this is not required.

**insert [OID = oid_value] (value1 value2 ...)**

Insert a new row into the open table using value1, value2, etc., for its column values and oid_value for its OID. If oid_value is zero (0) or the clause is omitted, and the table has OIDs, then the next available OID is assigned.

NULL values can be specified using the special key word _null_. Values containing spaces must be double quoted.

**declare [unique] index indexname indexoid on tablename using amname ( opclass1 name1 [, ...] )**

Create an index named indexname, having OID indexoid, on the table named tablename, using the amname access method. The fields to index are called name1, name2 etc., and the operator classes to use are opclass1, opclass2 etc., respectively. The index file is created and appropriate catalog entries are made for it, but the index contents are not initialized by this command.

**declare toast toasttableoid toastindexoid on tablename**

Create a TOAST table for the table named tablename. The TOAST table is assigned OID toasttableoid and its index is assigned OID toastindexoid. As with declare index, filling of the index is postponed.

**build indices**

Fill in the indices that have previously been declared.

### 66.3. Structure of the Bootstrap BKI File

The open command cannot be used until the tables it uses exist and have entries for the table that is to be opened. (These minimum tables are pg_class, pg_attribute, pg_proc, and pg_type.) To allow those tables themselves to be filled, create with the bootstrap option implicitly opens the created table for data insertion.

Also, the declare index and declare toast commands cannot be used until the system catalogs they need have been created and filled in.

Thus, the structure of the postgres.bki file has to be:

1. create bootstrap one of the critical tables
2. insert data describing at least the critical tables
3. close
4. Repeat for the other critical tables.
5. create (without bootstrap) a noncritical table
6. open
7. insert desired data
8. close
9. Repeat for the other noncritical tables.
10. Define indexes and toast tables.
11. build indices

There are doubtless other, undocumented ordering dependencies.

66.4. Example

The following sequence of commands will create the table test_table with OID 420, having two columns cola and colb of type int4 and text, respectively, and insert two rows into the table:

create test_table 420 (cola = int4, colb = text)
open test_table
insert OID=421 { 1 "value1" }
insert OID=422 { 2 _null_ }
close test_table
Chapter 67. How the Planner Uses Statistics

This chapter builds on the material covered in Section 14.1 and Section 14.2 to show some additional details about how the planner uses the system statistics to estimate the number of rows each part of a query might return. This is a significant part of the planning process, providing much of the raw material for cost calculation.

The intent of this chapter is not to document the code in detail, but to present an overview of how it works. This will perhaps ease the learning curve for someone who subsequently wishes to read the code.

67.1. Row Estimation Examples

The examples shown below use tables in the PostgreSQL regression test database. The outputs shown are taken from version 8.3. The behavior of earlier (or later) versions might vary. Note also that since ANALYZE uses random sampling while producing statistics, the results will change slightly after any new ANALYZE.

Let's start with a very simple query:

```sql
EXPLAIN SELECT * FROM tenk1;
```

```
QUERY PLAN

Seq Scan on tenk1 (cost=0.00..458.00 rows=10000 width=244)
```

How the planner determines the cardinality of tenk1 is covered in Section 14.2, but is repeated here for completeness. The number of pages and rows is looked up in pg_class:

```sql
SELECT relpages, reltuples FROM pg_class WHERE relname = 'tenk1';
```

```
relpages | reltuples
----------+-----------
358 | 10000
```

These numbers are current as of the last VACUUM or ANALYZE on the table. The planner then fetches the actual current number of pages in the table (this is a cheap operation, not requiring a table scan). If that is different from relpages then reltuples is scaled accordingly to arrive at a current number-of-rows estimate. In the example above, the value of relpages is up-to-date so the rows estimate is the same as reltuples.

Let's move on to an example with a range condition in its WHERE clause:

```sql
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 1000;
```

```
QUERY PLAN

Bitmap Heap Scan on tenk1 (cost=24.06..394.64 rows=1007 width=244)
  Recheck Cond: (unique1 < 1000)
  -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..23.80 rows=1007 width=0)
    Index Cond: (unique1 < 1000)
```
Chapter 67. How the Planner Uses Statistics

The planner examines the \texttt{WHERE} clause condition and looks up the selectivity function for the \texttt{operator <} in \texttt{pg\_operator}. This is held in the column \texttt{oprrest}, and the entry in this case is \texttt{scalarltsel}. The \texttt{scalarltsel} function retrieves the histogram for \texttt{unique1} from \texttt{pg\_statistics}. For manual queries it is more convenient to look in the simpler \texttt{pg\_stats} view:

```sql
SELECT histogram_bounds FROM pg_stats
WHERE table_name='tenk1' AND attname='unique1';
```

```
<table>
<thead>
<tr>
<th>histogram_bounds</th>
</tr>
</thead>
<tbody>
<tr>
<td>{0,993,1997,3050,4040,5036,5957,7057,8029,9016,9995}</td>
</tr>
</tbody>
</table>
```

Next the fraction of the histogram occupied by “< 1000” is worked out. This is the selectivity. The histogram divides the range into equal frequency buckets, so all we have to do is locate the bucket that our value is in and count part of it and all of the ones before. The value 1000 is clearly in the second bucket (993-1997). Assuming a linear distribution of values inside each bucket, we can calculate the selectivity as:

\[
\text{selectivity} = \left(1 + \frac{1000 - \text{bucket[2].min}}{\text{bucket[2].max} - \text{bucket[2].min}}\right)/\text{num\_buckets}
\]

\[
= \left(1 + \frac{1000 - 993}{1997 - 993}\right)/10
\]

\[
= 0.100697
\]

that is, one whole bucket plus a linear fraction of the second, divided by the number of buckets. The estimated number of rows can now be calculated as the product of the selectivity and the cardinality of \texttt{tenk1}:

\[
\text{rows} = \text{rel\_cardinality} \times \text{selectivity}
\]

\[
= 10000 \times 0.100697
\]

\[
= 1007 \text{ (rounding off)}
\]

Next let’s consider an example with an equality condition in its \texttt{WHERE} clause:

```sql
EXPLAIN SELECT * FROM tenk1 WHERE stringu1 = 'CRAAAA';
```

```
QUERY PLAN

Seq Scan on tenk1 (cost=0.00..483.00 rows=30 width=244)
Filter: (stringu1 = 'CRAAAA'::name)
```

Again the planner examines the \texttt{WHERE} clause condition and looks up the selectivity function for \texttt{=}, which is \texttt{eqsel}. For equality estimation the histogram is not useful; instead the list of most common values (MCVs) is used to determine the selectivity. Let’s have a look at the MCVs, with some additional columns that will be useful later:

```sql
SELECT null_frac, n_distinct, most_common_vals, most_common_freqs FROM pg_stats
WHERE table_name='tenk1' AND attname='stringu1';
```

```
| null_frac | 0 |
| n_distinct | 676 |
| most_common_vals | {EJAAAA, BBAAAA, CRAAAA, FCAAAA, FEAAAA, GSAAAA, JOAAAA, MCAAAA, NAAAAA, WGGAAA} |
| most_common_freqs | {0.00333333, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003, 0.003} |
```

Since \texttt{CRAAAA} appears in the list of MCVs, the selectivity is merely the corresponding entry in the list of most common frequencies (MCFs):
Chapter 67. How the Planner Uses Statistics

selectivity = mcf[3] = 0.003

As before, the estimated number of rows is just the product of this with the cardinality of tenk1:

rows = 10000 * 0.003 = 30

Now consider the same query, but with a constant that is not in the MCV list:

EXPLAIN SELECT * FROM tenk1 WHERE stringu1 = 'xxx';

QUERY PLAN
----------------------------------------------------------
Seq Scan on tenk1 (cost=0.00..483.00 rows=15 width=244)
Filter: (stringu1 = 'xxx'::name)

This is quite a different problem: how to estimate the selectivity when the value is not in the MCV list. The approach is to use the fact that the value is not in the list, combined with the knowledge of the frequencies for all of the MCVs:

selectivity = (1 - sum(mvf))/(num_distinct - num_mcv)
= (1 - (0.00333333 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003))/(676 - 10)
= 0.0014559

That is, add up all the frequencies for the MCVs and subtract them from one, then divide by the number of other distinct values. This amounts to assuming that the fraction of the column that is not any of the MCVs is evenly distributed among all the other distinct values. Notice that there are no null values so we don’t have to worry about those (otherwise we’d subtract the null fraction from the numerator as well). The estimated number of rows is then calculated as usual:

rows = 10000 * 0.0014559 = 15 (rounding off)

The previous example with unique1 < 1000 was an oversimplification of what scalarltsel really does; now that we have seen an example of the use of MCVs, we can fill in some more detail. The example was correct as far as it went, because since unique1 is a unique column it has no MCVs (obviously, no value is any more common than any other value). For a non-unique column, there will normally be both a histogram and an MCV list, and the histogram does not include the portion of the column population represented by the MCVs. We do things this way because it allows more precise estimation. In this situation scalarltsel directly applies the condition (e.g., "< 1000") to each value of the MCV list, and adds up the frequencies of the MCVs for which the condition is true. This gives an exact estimate of the selectivity within the portion of the table that is MCVs. The histogram is then used in the same way as above to estimate the selectivity in the portion of the table that is not MCVs, and then the two numbers are combined to estimate the overall selectivity. For example, consider

EXPLAIN SELECT * FROM tenk1 WHERE stringu1 < 'IAAAAA';

QUERY PLAN
------------------------------------------------------------
Seq Scan on tenk1 (cost=0.00..483.00 rows=3077 width=244)
Chapter 67. How the Planner Uses Statistics

Filter: (stringu1 < 'IAAAAA'::name)

We already saw the MCV information for stringu1, and here is its histogram:

```
SELECT histogram_bounds FROM pg_stats
WHERE tablename='tenk1' AND attname='stringu1';
```

```
histogram_bounds
---------------------------------------------
{AAAAAA, CQAAAA, FRAAAA, IBAAAA, KRAAAA, NFAAAA, PSAAAA, SGAAAA, VAAAAA, XLAAAA, ZZAAAA}
```

 Checking the MCV list, we find that the condition stringu1 < 'IAAAAA' is satisfied by the first six entries and not the last four, so the selectivity within the MCV part of the population is

\[
\text{selectivity} = \sum \text{relevant mvfs} = 0.00333333 + 0.003 + 0.003 + 0.003 + 0.003 + 0.003 = 0.01833333
\]

Summing all the MCFs also tells us that the total fraction of the population represented by MCVs is 0.03033333, and therefore the fraction represented by the histogram is 0.96966667 (again, there are no nulls, else we'd have to exclude them here). We can see that the value IAAAAA falls nearly at the end of the third histogram bucket. Using some rather cheesy assumptions about the frequency of different characters, the planner arrives at the estimate 0.298387 for the portion of the histogram population that is less than IAAAAA. We then combine the estimates for the MCV and non-MCV populations:

\[
\text{selectivity} = \text{mcv selectivity} + \text{histogram selectivity} \times \text{histogram fraction} = 0.01833333 + 0.298387 \times 0.96966667 = 0.307669
\]

\[
\text{rows} = 10000 \times 0.307669 = 3077 \text{ (rounding off)}
\]

In this particular example, the correction from the MCV list is fairly small, because the column distribution is actually quite flat (the statistics showing these particular values as being more common than others are mostly due to sampling error). In a more typical case where some values are significantly more common than others, this complicated process gives a useful improvement in accuracy because the selectivity for the most common values is found exactly.

Now let's consider a case with more than one condition in the WHERE clause:

```
EXPLAIN SELECT * FROM tenk1 WHERE unique1 < 1000 AND stringu1 = 'xxx';
```

```
QUERY PLAN
-------------------------------------------------------------------
Bitmap Heap Scan on tenk1 (cost=23.80..396.91 rows=1 width=244)  
  Recheck Cond: (unique1 < 1000)  
  Filter: (stringu1 = 'xxx'::name)  
    -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..23.80 rows=1007 width=0)  
      Index Cond: (unique1 < 1000)
```

The planner assumes that the two conditions are independent, so that the individual selectivities of the clauses can be multiplied together:

\[
\text{selectivity} = \text{selectivity(unique1 < 1000)} \times \text{selectivity(stringu1 = 'xxx')} = 0.100697 \times 0.0014559 = 0.0001466
\]
rows = 10000 \times 0.0001466
= 1 \text{ (rounding off)}

Notice that the number of rows estimated to be returned from the bitmap index scan reflects only the condition used with the index; this is important since it affects the cost estimate for the subsequent heap fetches.

Finally we will examine a query that involves a join:

```
EXPLAIN SELECT * FROM tenk1 t1, tenk2 t2
WHERE t1.unique1 < 50 AND t1.unique2 = t2.unique2;
```

```
QUERY PLAN
--------------------------------------------------------------------------------------
Nested Loop (cost=4.64..456.23 rows=50 width=488)
  -> Bitmap Heap Scan on tenk1 t1 (cost=4.64..142.17 rows=50 width=244)
      Recheck Cond: (unique1 < 50)
      -> Bitmap Index Scan on tenk1_unique1 (cost=0.00..4.63 rows=50 width=0)
          Index Cond: (unique1 < 50)
      -> Index Scan using tenk2_unique2 on tenk2 t2 (cost=0.00..6.27 rows=1 width=244)
          Index Cond: (unique2 = t1.unique2)
```

The restriction on \text{tenk1, unique1} < 50, is evaluated before the nested-loop join. This is handled analogously to the previous range example. This time the value 50 falls into the first bucket of the unique1 histogram:

\[
\text{selectivity} = \frac{(0 + (50 - \text{bucket}[1].\text{min})/(\text{bucket}[1].\text{max} - \text{bucket}[1].\text{min}))/\text{num_buckets}}{10} = 0.005035
\]

\[
\text{rows} = 10000 \times 0.005035
= 50 \text{ (rounding off)}
\]

The restriction for the join is \text{t2.unique2} = \text{t1.unique2}. The operator is just our familiar \text{=}, however the selectivity function is obtained from the \text{oprjoin} column of \text{pg_operator}, and is \text{eqjoinsel}. \text{eqjoinsel} looks up the statistical information for both \text{tenk2 and tenk1}:

```
SELECT tablename, null_frac,n_distinct, most_common_vals FROM pg_stats
WHERE tablename IN ('tenk1', 'tenk2') AND attname='unique2';
```

```
tablename | null_frac | n_distinct | most_common_vals
-----------------------------------------------
tenk1     | 0         | -1         | -1
ntenk2    | 0         | -1         | -1
```

In this case there is no MCV information for \text{unique2} because all the values appear to be unique, so we use an algorithm that relies only on the number of distinct values for both relations together with their null fractions:

\[
\text{selectivity} = (1 - \text{null_frac1}) \times (1 - \text{null_frac2}) \times \min(1/\text{n_distinct1}, 1/\text{n_distinct2})
= (1 - 0) \times (1 - 0) / \max(10000, 10000)
= 0.0001
\]

This is, subtract the null fraction from one for each of the relations, and divide by the maximum of the numbers of distinct values. The number of rows that the join is likely to emit is calculated as the cardinality of the Cartesian product of the two inputs, multiplied by the selectivity:
Chapter 67. How the Planner Uses Statistics

rows = (outer_cardinality * inner_cardinality) * selectivity
= (50 * 10000) * 0.0001
= 50

Had there been MCV lists for the two columns, `eqjoinsel` would have used direct comparison of the MCV lists to determine the join selectivity within the part of the column populations represented by the MCVs. The estimate for the remainder of the populations follows the same approach shown here.

Notice that we showed `inner_cardinality` as 10000, that is, the unmodified size of `tenk2`. It might appear from inspection of the `EXPLAIN` output that the estimate of join rows comes from 50 * 1, that is, the number of outer rows times the estimated number of rows obtained by each inner index scan on `tenk2`. But this is not the case: the join relation size is estimated before any particular join plan has been considered. If everything is working well then the two ways of estimating the join size will produce about the same answer, but due to round-off error and other factors they sometimes diverge significantly.

For those interested in further details, estimation of the size of a table (before any `WHERE` clauses) is done in `src/backend/optimizer/util/plancat.c`. The generic logic for clause selectivities is in `src/backend/optimizer/path/clausesel.c`. The operator-specific selectivity functions are mostly found in `src/backend/utils/adt/selfuncs.c`.

67.2. Planner Statistics and Security

Access to the table `pg_statistic` is restricted to superusers, so that ordinary users cannot learn about the contents of the tables of other users from it. Some selectivity estimation functions will use a user-provided operator (either the operator appearing in the query or a related operator) to analyze the stored statistics. For example, in order to determine whether a stored most common value is applicable, the selectivity estimator will have to run the appropriate operator to compare the constant in the query to the stored value. Thus the data in `pg_statistic` is potentially passed to user-defined operators. An appropriately crafted operator can intentionally leak the passed operands (for example, by logging them or writing them to a different table), or accidentally leak them by showing their values in error messages, in either case possibly exposing data from `pg_statistic` to a user who should not be able to see it.

In order to prevent this, the following applies to all built-in selectivity estimation functions. When planning a query, in order to be able to use stored statistics, the current user must either have `SELECT` privilege on the table or the involved columns, or the operator used must be `LEAKPROOF` (more accurately, the function that the operator is based on). If not, then the selectivity estimator will behave as if no statistics are available, and the planner will proceed with default or fall-back assumptions.

If a user does not have the required privilege on the table or columns, then in many cases the query will ultimately receive a permission-denied error, in which case this mechanism is invisible in practice. But if the user is reading from a security-barrier view, then the planner might wish to check the statistics of an underlying table that is otherwise inaccessible to the user. In that case, the operator should be leak-proof or the statistics will not be used. There is no direct feedback about that, except that the plan might be suboptimal. If one suspects that this is the case, one could try running the query as a more privileged user, to see if a different plan results.

This restriction applies only to cases where the planner would need to execute a user-defined operator on one or more values from `pg_statistic`. Thus the planner is permitted to use generic statistical...
information, such as the fraction of null values or the number of distinct values in a column, regardless of access privileges.

Selectivity estimation functions contained in third-party extensions that potentially operate on statistics with user-defined operators should follow the same security rules. Consult the PostgreSQL source code for guidance.
VIII. Appendixes
Appendix A. PostgreSQL Error Codes

All messages emitted by the PostgreSQL server are assigned five-character error codes that follow the SQL standard’s conventions for “SQLSTATE” codes. Applications that need to know which error condition has occurred should usually test the error code, rather than looking at the textual error message. The error codes are less likely to change across PostgreSQL releases, and also are not subject to change due to localization of error messages. Note that some, but not all, of the error codes produced by PostgreSQL are defined by the SQL standard; some additional error codes for conditions not defined by the standard have been invented or borrowed from other databases.

According to the standard, the first two characters of an error code denote a class of errors, while the last three characters indicate a specific condition within that class. Thus, an application that does not recognize the specific error code might still be able to infer what to do from the error class.

Table A-1 lists all the error codes defined in PostgreSQL 9.6.17. (Some are not actually used at present, but are defined by the SQL standard.) The error classes are also shown. For each error class there is a “standard” error code having the last three characters 000. This code is used only for error conditions that fall within the class but do not have any more-specific code assigned.

The symbol shown in the column “Condition Name” is the condition name to use in PL/pgSQL. Condition names can be written in either upper or lower case. (Note that PL/pgSQL does not recognize warning, as opposed to error, condition names; those are classes 00, 01, and 02.)

For some types of errors, the server reports the name of a database object (a table, table column, data type, or constraint) associated with the error; for example, the name of the unique constraint that caused a unique_violation error. Such names are supplied in separate fields of the error report message so that applications need not try to extract them from the possibly-localized human-readable text of the message. As of PostgreSQL 9.3, complete coverage for this feature exists only for errors in SQLSTATE class 23 (integrity constraint violation), but this is likely to be expanded in future.

Table A-1. PostgreSQL Error Codes

<table>
<thead>
<tr>
<th>Error Code</th>
<th>Condition Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>00000</td>
<td>successful_completion</td>
</tr>
<tr>
<td>Class 01 — Warning</td>
<td></td>
</tr>
<tr>
<td>01000</td>
<td>warning</td>
</tr>
<tr>
<td>0100C</td>
<td>dynamic_result_sets_returned</td>
</tr>
<tr>
<td>01008</td>
<td>implicit_zero_bit_padding</td>
</tr>
<tr>
<td>01003</td>
<td>null_value_eliminated_in_set_function</td>
</tr>
<tr>
<td>01007</td>
<td>privilege_not_granted</td>
</tr>
<tr>
<td>01006</td>
<td>privilege_not_revoked</td>
</tr>
<tr>
<td>01004</td>
<td>string_data_right_truncation</td>
</tr>
<tr>
<td>01P01</td>
<td>deprecated_feature</td>
</tr>
<tr>
<td>Class 02 — No Data (this is also a warning class per the SQL standard)</td>
<td></td>
</tr>
<tr>
<td>02000</td>
<td>no_data</td>
</tr>
<tr>
<td>02001</td>
<td>no_additional_dynamic_result_sets_returned</td>
</tr>
<tr>
<td>Class 03 — SQL Statement Not Yet Complete</td>
<td></td>
</tr>
<tr>
<td>Error Code</td>
<td>Condition Name</td>
</tr>
<tr>
<td>------------</td>
<td>----------------</td>
</tr>
<tr>
<td>03000</td>
<td>sql_statement_not_yet_complete</td>
</tr>
<tr>
<td></td>
<td><strong>Class 08 — Connection Exception</strong></td>
</tr>
<tr>
<td>08000</td>
<td>connection_exception</td>
</tr>
<tr>
<td>08003</td>
<td>connection_does_not_exist</td>
</tr>
<tr>
<td>08006</td>
<td>connection_failure</td>
</tr>
<tr>
<td>08001</td>
<td>sqlclient_unable_to_establish_sqlconnection</td>
</tr>
<tr>
<td>08004</td>
<td>sqlserver_rejected_establishment_of_sqlconnection</td>
</tr>
<tr>
<td>08007</td>
<td>transaction_resolution_unknown</td>
</tr>
<tr>
<td>08P01</td>
<td>protocolViolation</td>
</tr>
<tr>
<td></td>
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Appendix A. PostgreSQL Error Codes
### Appendix A. PostgreSQL Error Codes

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#### Class 23 — Integrity Constraint Violation

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#### Class 25 — Invalid Transaction State

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#### Class 26 — Invalid SQL Statement Name

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#### Class 27 — Triggered Data Change Violation

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#### Class 28 — Invalid Authorization Specification

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<tr>
<td>HV00C</td>
<td>fdw_invalid_option_index</td>
</tr>
<tr>
<td>HV00D</td>
<td>fdw_invalid_option_name</td>
</tr>
<tr>
<td>HV090</td>
<td>fdw_invalid_string_length_or_buffer_length</td>
</tr>
<tr>
<td>HV00A</td>
<td>fdw_invalid_string_format</td>
</tr>
<tr>
<td>HV009</td>
<td>fdw_invalid_use_of_null_pointer</td>
</tr>
<tr>
<td>HV014</td>
<td>fdw_too_many_handles</td>
</tr>
<tr>
<td>HV001</td>
<td>fdw_out_of_memory</td>
</tr>
<tr>
<td>HV00P</td>
<td>fdw_no_schemas</td>
</tr>
<tr>
<td>HV00J</td>
<td>fdw_option_name_not_found</td>
</tr>
<tr>
<td>HV00K</td>
<td>fdw_reply_handle</td>
</tr>
<tr>
<td>HV00Q</td>
<td>fdw_schema_not_found</td>
</tr>
<tr>
<td>HV00R</td>
<td>fdw_table_not_found</td>
</tr>
<tr>
<td>HV00L</td>
<td>fdw_unable_to_create_execution</td>
</tr>
<tr>
<td>HV00M</td>
<td>fdw_unable_to_create_reply</td>
</tr>
<tr>
<td>HV00N</td>
<td>fdw_unable_to_establish_connection</td>
</tr>
</tbody>
</table>

**Class P0 — PL/pgSQL Error**

| P0000      | plpgsql_error                                      |
| P0001      | raise_exception                                    |
| P0002      | no_data_found                                      |
| P0003      | too_many_rows                                      |
| P0004      | assert_failure                                     |

**Class XX — Internal Error**

| XX000      | internal_error                                     |
| XX001      | data_corrupted                                     |
| XX002      | index_corrupted                                     |
Appendix B. Date/Time Support

PostgreSQL uses an internal heuristic parser for all date/time input support. Dates and times are input as strings, and are broken up into distinct fields with a preliminary determination of what kind of information can be in the field. Each field is interpreted and either assigned a numeric value, ignored, or rejected. The parser contains internal lookup tables for all textual fields, including months, days of the week, and time zones.

This appendix includes information on the content of these lookup tables and describes the steps used by the parser to decode dates and times.

B.1. Date/Time Input Interpretation

Date/time input strings are decoded using the following procedure.

1. Break the input string into tokens and categorize each token as a string, time, time zone, or number.
   a. If the numeric token contains a colon (:) this is a time string. Include all subsequent digits and colons.
   b. If the numeric token contains a dash (-), slash (/), or two or more dots (.), this is a date string which might have a text month. If a date token has already been seen, it is instead interpreted as a time zone name (e.g., America/New_York).
   c. If the token is numeric only, then it is either a single field or an ISO 8601 concatenated date (e.g., 19990113 for January 13, 1999) or time (e.g., 141516 for 14:15:16).
   d. If the token starts with a plus (+) or minus (-), then it is either a numeric time zone or a special field.

2. If the token is an alphabetic string, match up with possible strings:
   a. See if the token matches any known time zone abbreviation. These abbreviations are supplied by the configuration file described in Section B.4.
   b. If not found, search an internal table to match the token as either a special string (e.g., today), day (e.g., Thursday), month (e.g., January), or noise word (e.g., at, on).
   c. If still not found, throw an error.

3. When the token is a number or number field:
   a. If there are eight or six digits, and if no other date fields have been previously read, then interpret as a “concatenated date” (e.g., 19990118 or 990118). The interpretation is YYYYMMDD or YYMMDD.
   b. If the token is three digits and a year has already been read, then interpret as day of year.
   c. If four or six digits and a year has already been read, then interpret as a time (HHMM or HHMMSS).

2210
Appendix B. Date/Time Support

d. If three or more digits and no date fields have yet been found, interpret as a year (this forces yy-mm-dd ordering of the remaining date fields).
e. Otherwise the date field ordering is assumed to follow the DateStyle setting: mm-dd-yy, dd-mm-yy, or yy-mm-dd. Throw an error if a month or day field is found to be out of range.

4. If BC has been specified, negate the year and add one for internal storage. (There is no year zero in the Gregorian calendar, so numerically 1 BC becomes year zero.)

5. If BC was not specified, and if the year field was two digits in length, then adjust the year to four digits. If the field is less than 70, then add 2000, otherwise add 1900.

Tip: Gregorian years AD 1-99 can be entered by using 4 digits with leading zeros (e.g., 0099 is AD 99).

B.2. Handling of Invalid or Ambiguous Timestamps

Ordinarily, if a date/time string is syntactically valid but contains out-of-range field values, an error will be thrown. For example, input specifying the 31st of February will be rejected.

During a daylight-savings-time transition, it is possible for a seemingly valid timestamp string to represent a nonexistent or ambiguous timestamp. Such cases are not rejected; the ambiguity is resolved by determining which UTC offset to apply. For example, supposing that the TimeZone parameter is set to America/New_York, consider

```sql
=> SELECT '2018-03-11 02:30'::timestamptz;
timestamptz
------------------------
2018-03-11 03:30:00-04 (1 row)
```

Because that day was a spring-forward transition date in that time zone, there was no civil time instant 2:30AM; clocks jumped forward from 2AM EST to 3AM EDT. PostgreSQL interprets the given time as if it were standard time (UTC-5), which then renders as 3:30AM EDT (UTC-4).

Conversely, consider the behavior during a fall-back transition:

```sql
=> SELECT '2018-11-04 02:30'::timestamptz;
timestamptz
------------------------
2018-11-04 02:30:00-05 (1 row)
```

On that date, there were two possible interpretations of 2:30AM; there was 2:30AM EDT, and then an hour later after the reversion to standard time, there was 2:30AM EST. Again, PostgreSQL interprets the given time as if it were standard time (UTC-5). We can force the matter by specifying daylight-savings time:

```sql
=> SELECT '2018-11-04 02:30 EDT'::timestamptz;
timestamptz
------------------------
```

2211
This timestamp could validly be rendered as either 2:30 UTC-4 or 1:30 UTC-5; the timestamp output code chooses the latter.

The precise rule that is applied in such cases is that an invalid timestamp that appears to fall within a jump-forward daylight savings transition is assigned the UTC offset that prevailed in the time zone just before the transition, while an ambiguous timestamp that could fall on either side of a jump-back transition is assigned the UTC offset that prevailed just after the transition. In most time zones this is equivalent to saying that “the standard-time interpretation is preferred when in doubt”.

In all cases, the UTC offset associated with a timestamp can be specified explicitly, using either a numeric UTC offset or a time zone abbreviation that corresponds to a fixed UTC offset. The rule just given applies only when it is necessary to infer a UTC offset for a time zone in which the offset varies.

### B.3. Date/Time Key Words

Table B-1 shows the tokens that are recognized as names of months.

**Table B-1. Month Names**

<table>
<thead>
<tr>
<th>Month</th>
<th>Abbreviations</th>
</tr>
</thead>
<tbody>
<tr>
<td>January</td>
<td>Jan</td>
</tr>
<tr>
<td>February</td>
<td>Feb</td>
</tr>
<tr>
<td>March</td>
<td>Mar</td>
</tr>
<tr>
<td>April</td>
<td>Apr</td>
</tr>
<tr>
<td>May</td>
<td></td>
</tr>
<tr>
<td>June</td>
<td>Jun</td>
</tr>
<tr>
<td>July</td>
<td>Jul</td>
</tr>
<tr>
<td>August</td>
<td>Aug</td>
</tr>
<tr>
<td>September</td>
<td>Sep, Sept</td>
</tr>
<tr>
<td>October</td>
<td>Oct</td>
</tr>
<tr>
<td>November</td>
<td>Nov</td>
</tr>
<tr>
<td>December</td>
<td>Dec</td>
</tr>
</tbody>
</table>

Table B-2 shows the tokens that are recognized as names of days of the week.

**Table B-2. Day of the Week Names**

<table>
<thead>
<tr>
<th>Day</th>
<th>Abbreviations</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sunday</td>
<td>Sun</td>
</tr>
<tr>
<td>Monday</td>
<td>Mon</td>
</tr>
<tr>
<td>Tuesday</td>
<td>Tue, Tues</td>
</tr>
<tr>
<td>Wednesday</td>
<td>Wed, Weds</td>
</tr>
<tr>
<td>Thursday</td>
<td>Thu, Thur, Thurs</td>
</tr>
<tr>
<td>Friday</td>
<td>Fri</td>
</tr>
<tr>
<td>Saturday</td>
<td>Sat</td>
</tr>
</tbody>
</table>
Table B-3 shows the tokens that serve various modifier purposes.

<table>
<thead>
<tr>
<th>Identifier</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AM</td>
<td>Time is before 12:00</td>
</tr>
<tr>
<td>AT</td>
<td>Ignored</td>
</tr>
<tr>
<td>JULIAN, JD, J</td>
<td>Next field is Julian Date</td>
</tr>
<tr>
<td>ON</td>
<td>Ignored</td>
</tr>
<tr>
<td>PM</td>
<td>Time is on or after 12:00</td>
</tr>
<tr>
<td>T</td>
<td>Next field is time</td>
</tr>
</tbody>
</table>

### B.4. Date/Time Configuration Files

Since timezone abbreviations are not well standardized, PostgreSQL provides a means to customize the set of abbreviations accepted by the server. The timezone_abbreviations run-time parameter determines the active set of abbreviations. While this parameter can be altered by any database user, the possible values for it are under the control of the database administrator — they are in fact names of configuration files stored in `.../share/timezonesets/` of the installation directory. By adding or altering files in that directory, the administrator can set local policy for timezone abbreviations.

The `timezone_abbreviations` can be set to any file name found in `.../share/timezonesets/`, if the file’s name is entirely alphabetic. (The prohibition against non-alphabetic characters in `timezone_abbreviations` prevents reading files outside the intended directory, as well as reading editor backup files and other extraneous files.)

A timezone abbreviation file can contain blank lines and comments beginning with `#`. Non-comment lines must have one of these formats:

- `zone_abbreviation offset`
- `zone_abbreviation offset D`
- `zone_abbreviation time_zone_name`
- `@INCLUDE file_name`
- `@OVERRIDE`

A `zone_abbreviation` is just the abbreviation being defined. An `offset` is an integer giving the equivalent offset in seconds from UTC, positive being east from Greenwich and negative being west. For example, -18000 would be five hours west of Greenwich, or North American east coast standard time. `D` indicates that the zone name represents local daylight-savings time rather than standard time.

Alternatively, a `time_zone_name` can be given, referencing a zone name defined in the IANA timezone database. The zone’s definition is consulted to see whether the abbreviation is or has been in use in that zone, and if so, the appropriate meaning is used — that is, the meaning that was currently in use at the timestamp whose value is being determined, or the meaning in use immediately before that if it wasn’t current at that time, or the oldest meaning if it was used only after that time. This behavior is essential for dealing with abbreviations whose meaning has historically varied. It is also allowed to define an abbreviation in terms of a zone name in which that abbreviation does not appear; then using the abbreviation is just equivalent to writing out the zone name.
Appendix B. Date/Time Support

**Tip:** Using a simple integer *offset* is preferred when defining an abbreviation whose offset from UTC has never changed, as such abbreviations are much cheaper to process than those that require consulting a time zone definition.

The `@INCLUDE` syntax allows inclusion of another file in the `.../share/timezonesets/` directory. Inclusion can be nested, to a limited depth.

The `@OVERRIDE` syntax indicates that subsequent entries in the file can override previous entries (typically, entries obtained from included files). Without this, conflicting definitions of the same timezone abbreviation are considered an error.

In an unmodified installation, the file `Default` contains all the non-conflicting time zone abbreviations for most of the world. Additional files `Australia` and `India` are provided for those regions: these files first include the `Default` file and then add or modify abbreviations as needed.

For reference purposes, a standard installation also contains files `Africa.txt`, `America.txt`, etc, containing information about every time zone abbreviation known to be in use according to the IANA timezone database. The zone name definitions found in these files can be copied and pasted into a custom configuration file as needed. Note that these files cannot be directly referenced as `timezone_abbreviations` settings, because of the dot embedded in their names.

**Note:** If an error occurs while reading the time zone abbreviation set, no new value is applied and the old set is kept. If the error occurs while starting the database, startup fails.

---

**Caution**

Time zone abbreviations defined in the configuration file override non-timezone meanings built into PostgreSQL. For example, the `Australia` configuration file defines `SAT` (for South Australian Standard Time). When this file is active, `SAT` will not be recognized as an abbreviation for Saturday.

---

**Caution**

If you modify files in `.../share/timezonesets/`, it is up to you to make backups — a normal database dump will not include this directory.

---

**B.5. History of Units**

The SQL standard states that “Within the definition of a ‘datetime literal’, the ‘datetime values’ are constrained by the natural rules for dates and times according to the Gregorian calendar”. PostgreSQL follows the SQL standard’s lead by counting dates exclusively in the Gregorian calendar, even for years before that calendar was in use. This rule is known as the *proleptic* Gregorian calendar.

The Julian calendar was introduced by Julius Caesar in 45 BC. It was in common use in the Western world until the year 1582, when countries started changing to the Gregorian calendar. In the Julian calendar, the tropical year is approximated as 365 1/4 days = 365.25 days. This gives an error of about 1 day in 128 years.

The accumulating calendar error prompted Pope Gregory XIII to reform the calendar in accordance with instructions from the Council of Trent. In the Gregorian calendar, the tropical year is approxi-
mated as 365 + 97 / 400 days = 365.2425 days. Thus it takes approximately 3300 years for the tropical year to shift one day with respect to the Gregorian calendar.

The approximation 365+97/400 is achieved by having 97 leap years every 400 years, using the following rules:

Every year divisible by 4 is a leap year.
However, every year divisible by 100 is not a leap year.
However, every year divisible by 400 is a leap year after all.

So, 1700, 1800, 1900, 2100, and 2200 are not leap years. But 1600, 2000, and 2400 are leap years. By contrast, in the older Julian calendar all years divisible by 4 are leap years.

The papal bull of February 1582 decreed that 10 days should be dropped from October 1582 so that 15 October should follow immediately after 4 October. This was observed in Italy, Poland, Portugal, and Spain. Other Catholic countries followed shortly after, but Protestant countries were reluctant to change, and the Greek Orthodox countries didn’t change until the start of the 20th century. The reform was observed by Great Britain and its dominions (including what is now the USA) in 1752. Thus 2 September 1752 was followed by 14 September 1752. This is why Unix systems that have the cal program produce the following:

```
$ cal 9 1752
  September 1752
 Su Mo Tu We Th Fr Sa
  1   14  15  16
 2 17 18 19 20 21 22 23
24 25 26 27 28 29 30
```

But, of course, this calendar is only valid for Great Britain and dominions, not other places. Since it would be difficult and confusing to try to track the actual calendars that were in use in various places at various times, PostgreSQL does not try, but rather follows the Gregorian calendar rules for all dates, even though this method is not historically accurate.

Different calendars have been developed in various parts of the world, many predating the Gregorian system. For example, the beginnings of the Chinese calendar can be traced back to the 14th century BC. Legend has it that the Emperor Huangdi invented that calendar in 2637 BC. The People’s Republic of China uses the Gregorian calendar for civil purposes. The Chinese calendar is used for determining festivals.

The Julian Date system is another type of calendar, unrelated to the Julian calendar though it is confusingly named similarly to that calendar. The Julian Date system was invented by the French scholar Joseph Justus Scaliger (1540-1609) and probably takes its name from Scaliger’s father, the Italian scholar Julius Caesar Scaliger (1484-1558). In the Julian Date system, each day has a sequential number, starting from JD 0 (which is sometimes called the Julian Date). JD 0 corresponds to 1 January 4713 BC in the Julian calendar, or 24 November 4714 BC in the Gregorian calendar. Julian Date counting is most often used by astronomers for labeling their nightly observations, and therefore a date runs from noon UTC to the next noon UTC, rather than from midnight to midnight: JD 0 designates the 24 hours from noon UTC on 24 November 4714 BC to noon UTC on 25 November 4714 BC.

Although PostgreSQL supports Julian Date notation for input and output of dates (and also uses Julian dates for some internal datetime calculations), it does not observe the nicety of having dates run from noon to noon. PostgreSQL treats a Julian Date as running from midnight to midnight.
Appendix C. SQL Key Words

Table C-1 lists all tokens that are key words in the SQL standard and in PostgreSQL 9.6.17. Background information can be found in Section 4.1.1. (For space reasons, only the latest two versions of the SQL standard, and SQL-92 for historical comparison, are included. The differences between those and the other intermediate standard versions are small.)

SQL distinguishes between reserved and non-reserved key words. According to the standard, reserved key words are the only real key words; they are never allowed as identifiers. Non-reserved key words only have a special meaning in particular contexts and can be used as identifiers in other contexts. Most non-reserved key words are actually the names of built-in tables and functions specified by SQL. The concept of non-reserved key words essentially only exists to declare that some predefined meaning is attached to a word in some contexts.

In the PostgreSQL parser life is a bit more complicated. There are several different classes of tokens ranging from those that can never be used as an identifier to those that have absolutely no special status in the parser as compared to an ordinary identifier. (The latter is usually the case for functions specified by SQL.) Even reserved key words are not completely reserved in PostgreSQL, but can be used as column labels (for example, SELECT 55 AS CHECK, even though CHECK is a reserved key word).

In Table C-1 in the column for PostgreSQL we classify as “non-reserved” those key words that are explicitly known to the parser but are allowed as column or table names. Some key words that are otherwise non-reserved cannot be used as function or data type names and are marked accordingly. (Most of these words represent built-in functions or data types with special syntax. The function or type is still available but it cannot be redefined by the user.) Labeled “reserved” are those tokens that are not allowed as column or table names. Some reserved key words are allowable as names for functions or data types; this is also shown in the table. If not so marked, a reserved key word is only allowed as an “AS” column label name.

As a general rule, if you get spurious parser errors for commands that contain any of the listed key words as an identifier you should try to quote the identifier to see if the problem goes away.

It is important to understand before studying Table C-1 that the fact that a key word is not reserved in PostgreSQL does not mean that the feature related to the word is not implemented. Conversely, the presence of a key word does not indicate the existence of a feature.

Table C-1. SQL Key Words

<table>
<thead>
<tr>
<th>Key Word</th>
<th>PostgreSQL</th>
<th>SQL:2011</th>
<th>SQL:2008</th>
<th>SQL-92</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ABORT</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ABS</td>
<td>reserved</td>
<td>reserved</td>
<td></td>
<td></td>
</tr>
<tr>
<td>ABSENT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td></td>
</tr>
<tr>
<td>ABSOLUTE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ACCESS</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ACCORDING</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ACTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ADA</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>Key Word</td>
<td>PostgreSQL</td>
<td>SQL:2011</td>
<td>SQL:2008</td>
<td>SQL-92</td>
</tr>
<tr>
<td>---------------</td>
<td>------------</td>
<td>----------</td>
<td>----------</td>
<td>--------</td>
</tr>
<tr>
<td>ADD</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ADMIN</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>AFTER</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>AGGREGATE</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALL</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ALLOCATE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ALSO</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>ALTER</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ALWAYS</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ANALYZE</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>AND</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ANY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ARE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ARRAY</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ARRAY_AGG</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ARRAY_MAX_CARDINALITY</td>
<td>reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>AS</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASC</td>
<td>reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASSENSITIVE</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASSERTION</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ASSIGNMENT</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ASYMMETRIC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>AT</td>
<td>non-reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ATOMIC</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>ATTRIBUTE</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>ATTRIBUTES</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
<td>non-reserved</td>
</tr>
<tr>
<td>AUTHORIZATION</td>
<td>reserved (can be function or type)</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>AVG</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
<td>reserved</td>
</tr>
<tr>
<td>BACKWARD</td>
<td>non-reserved</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
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# Appendix C. SQL Key Words

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### Appendix C. SQL Key Words

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Appendix D. SQL Conformance

This section attempts to outline to what extent PostgreSQL conforms to the current SQL standard. The following information is not a full statement of conformance, but it presents the main topics in as much detail as is both reasonable and useful for users.

The formal name of the SQL standard is ISO/IEC 9075 “Database Language SQL”. A revised version of the standard is released from time to time; the most recent update appearing in 2011. The 2011 version is referred to as ISO/IEC 9075:2011, or simply as SQL:2011. The versions prior to that were SQL:2008, SQL:2003, SQL:1999, and SQL-92. Each version replaces the previous one, so claims of conformance to earlier versions have no official merit. PostgreSQL development aims for conformance with the latest official version of the standard where such conformance does not contradict traditional features or common sense. Many of the features required by the SQL standard are supported, though sometimes with slightly differing syntax or function. Further moves towards conformance can be expected over time.

SQL-92 defined three feature sets for conformance: Entry, Intermediate, and Full. Most database management systems claiming SQL standard conformance were conforming at only the Entry level, since the entire set of features in the Intermediate and Full levels was either too voluminous or in conflict with legacy behaviors.

Starting with SQL:1999, the SQL standard defines a large set of individual features rather than the ineffectively broad three levels found in SQL-92. A large subset of these features represents the “Core” features, which every conforming SQL implementation must supply. The rest of the features are purely optional. Some optional features are grouped together to form “packages”, which SQL implementations can claim conformance to, thus claiming conformance to particular groups of features.

The standard versions beginning with SQL:2003 are also split into a number of parts. Each is known by a shorthand name. Note that these parts are not consecutively numbered.

- ISO/IEC 9075-1 Framework (SQL/Framework)
- ISO/IEC 9075-2 Foundation (SQL/Foundation)
- ISO/IEC 9075-3 Call Level Interface (SQL/CLI)
- ISO/IEC 9075-4 Persistent Stored Modules (SQL/PSM)
- ISO/IEC 9075-9 Management of External Data (SQL/MED)
- ISO/IEC 9075-10 Object Language Bindings (SQL/OLB)
- ISO/IEC 9075-11 Information and Definition Schemas (SQL/Schemata)
- ISO/IEC 9075-13 Routines and Types using the Java Language (SQL/JRT)
- ISO/IEC 9075-14 XML-related specifications (SQL/XML)

The PostgreSQL core covers parts 1, 2, 9, 11, and 14. Part 3 is covered by the ODBC driver, and part 13 is covered by the PL/Java plug-in, but exact conformance is currently not being verified for these components. There are currently no implementations of parts 4 and 10 for PostgreSQL.

PostgreSQL supports most of the major features of SQL:2011. Out of 179 mandatory features required for full Core conformance, PostgreSQL conforms to at least 160. In addition, there is a long list of
supported optional features. It might be worth noting that at the time of writing, no current version of any database management system claims full conformance to Core SQL:2011.

In the following two sections, we provide a list of those features that PostgreSQL supports, followed by a list of the features defined in SQL:2011 which are not yet supported in PostgreSQL. Both of these lists are approximate: There might be minor details that are nonconforming for a feature that is listed as supported, and large parts of an unsupported feature might in fact be implemented. The main body of the documentation always contains the most accurate information about what does and does not work.

**Note:** Feature codes containing a hyphen are subfeatures. Therefore, if a particular subfeature is not supported, the main feature is listed as unsupported even if some other subfeatures are supported.

### D.1. Supported Features

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<td>All comparison operators are supported (rather than just =)</td>
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<td>TIME data type (including support of TIME literal) with fractional seconds precision of at least 0</td>
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### D.2. Unsupported Features

The following features defined in SQL:2011 are not implemented in this release of PostgreSQL. In a few cases, equivalent functionality is available.

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<td>XMLTable: ordinality column option</td>
<td></td>
</tr>
<tr>
<td>X303</td>
<td></td>
<td>XMLTable: column default option</td>
<td></td>
</tr>
<tr>
<td>X304</td>
<td></td>
<td>XMLTable: passing a context item</td>
<td></td>
</tr>
<tr>
<td>X305</td>
<td></td>
<td>XMLTable: initializing an XQuery variable</td>
<td></td>
</tr>
</tbody>
</table>
Appendix E. Release Notes

The release notes contain the significant changes in each PostgreSQL release, with major features and migration issues listed at the top. The release notes do not contain changes that affect only a few users or changes that are internal and therefore not user-visible. For example, the optimizer is improved in almost every release, but the improvements are usually observed by users as simply faster queries.

A complete list of changes for each release can be obtained by viewing the Git logs for each release. The `pgsql-committers` email list¹ records all source code changes as well. There is also a web interface² that shows changes to specific files.

The name appearing next to each item represents the major developer for that item. Of course all changes involve community discussion and patch review, so each item is truly a community effort.

E.1. Release 9.6.17

Release date: 2020-02-13

This release contains a variety of fixes from 9.6.16. For information about new features in the 9.6 major release, see Section E.18.

E.1.1. Migration to Version 9.6.17

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.16, see Section E.2.

E.1.2. Changes

- Add missing permissions checks for `ALTER ... DEPENDS ON EXTENSION` (Álvaro Herrera)
  
  Marking an object as dependent on an extension did not have any privilege check whatsoever. This oversight allowed any user to mark routines, triggers, materialized views, or indexes as droppable by anyone able to drop an extension. Require that the calling user own the specified object (and hence have privilege to drop it). (CVE-2020-1720)

- Avoid failure in logical decoding when a large transaction must be spilled into many separate temporary files (Amit Khandekar)

- Fix failure in logical replication publisher after a database crash and restart (Vignesh C)

- Prevent premature shutdown of a Gather or GatherMerge plan node that is underneath a Limit node (Amit Kapila)
  
  This avoids failure if such a plan node needs to be scanned more than once, as for instance if it is on the inside of a nestloop.

- Avoid memory leak when there are no free dynamic shared memory slots (Thomas Munro)

¹. https://archives.postgresql.org/pgsql-committers/
². https://git.postgresql.org/gitweb?p=postgresql.git;a=summary
Appendix E. Release Notes

- Ignore the `CONCURRENTLY` option when performing an index creation, drop, or rebuild on a temporary table (Michael Paquier, Heikki Linnakangas, Andres Freund)
  
  This avoids strange failures if the temporary table has an `ON COMMIT` action. There is no benefit in using `CONCURRENTLY` for a temporary table anyway, since other sessions cannot access the table, making the extra processing pointless.

- Fix possible failure when resetting expression indexes on temporary tables that are marked `ON COMMIT DELETE ROWS` (Tom Lane)

- Fix possible crash in BRIN index operations with `box`, `range` and `inet` data types (Heikki Linnakangas)

- Fix handling of deleted pages in GIN indexes (Alexander Korotkov)
  
  Avoid possible deadlocks, incorrect updates of a deleted page’s state, and failure to traverse through a recently-deleted page.

- Fix possible crash with a SubPlan (sub-`SELECT`) within a multi-row `VALUES` list (Tom Lane)

- Fix unlikely crash with pass-by-reference aggregate transition states (Andres Freund, Teodor Sigaev)

- Improve error reporting in `to_date()` and `to_timestamp()` (Tom Lane, Álvaro Herrera)
  
  Reports about incorrect month or day names in input strings could truncate the input in the middle of a multi-byte character, leading to an improperly encoded error message that could cause follow-on failures. Truncate at the next whitespace instead.

- Fix off-by-one result for `EXTRACT(ISOYEAR FROM timestamp)` for BC dates (Tom Lane)

- Avoid stack overflow in `information_schema` views when a self-referential view exists in the system catalogs (Tom Lane)
  
  A self-referential view can’t work; it will always result in infinite recursion. We handled that situation correctly when trying to execute the view, but not when inquiring whether it is automatically updatable.

- Improve performance of hash joins with very large inner relations (Thomas Munro)

- Fix edge-case crashes and misestimations in selectivity calculations for the `<@` and `@>` range operators (Michael Paquier, Andrey Borodin, Tom Lane)

- Improve error reporting for attempts to use automatic updating of views with conditional `INSTEAD` rules (Dean Rasheed)
  
  This has never been supported, but previously the error was thrown only at execution time, so that it could be masked by planner errors.

- Prevent a composite type from being included in itself indirectly via a range type (Tom Lane, Julien Rouhau)

- Fix error reporting for index expressions of prohibited types (Amit Langote)

- Fix dumping of views that contain only a `VALUES` list to handle cases where a view output column has been renamed (Tom Lane)

- Transmit incoming `NOTIFY` messages to the client before sending `ReadyForQuery`, rather than after (Tom Lane)
  
  This change ensures that, with libpq and other client libraries that act similarly to it, any notifications received during a transaction will be available by the time the client thinks the transaction is complete. This probably makes no difference in practical applications (which would need to
cope with asynchronous notifications in any case); but it makes it easier to build test cases with reproducible behavior.

- Allow libpq to parse all GSS-related connection parameters even when the GSSAPI code hasn’t been compiled in (Tom Lane)
  
  This makes the behavior similar to our SSL support, where it was long ago deemed to be a good idea to always accept all the related parameters, even if some are ignored or restricted due to lack of the feature in a particular build.

- Fix incorrect handling of %b and %B format codes in ecpg’s PGTYPEStimestamp_fmt_asc() function (Tomas Vondra)
  
  Due to an off-by-one error, these codes would print the wrong month name, or possibly crash.

- Fix parallel pg_dump/pg_restore to more gracefully handle failure to create worker processes (Tom Lane)

- Prevent possible crash or lockup when attempting to terminate a parallel pg_dump/pg_restore run via a signal (Tom Lane)

- In pg_upgrade, look inside arrays and ranges while searching for non-upgradable data types in tables (Tom Lane)

- Apply more thorough syntax checking to createuser’s --connection-limit option (Álvaro Herrera)

- Avoid crash in postgres_fdw when trying to send a command like UPDATE remote_tab SET (x,y) = (SELECT ...) to the remote server (Tom Lane)

- In contrib/dict_int, reject maxlen settings less than one (Tomas Vondra)
  
  This prevents a possible crash with silly settings for that parameter.

-Disallow NULL category values in contrib/tablefunc's crosstab() function (Joe Conway)
  
  This case never worked usefully, and it would crash on some platforms.

- Mark some timeout and statistics-tracking GUC variables as PGDLLIMPORT, to allow extensions to access them on Windows (Pascal Legrand)
  
  This applies to idle_in_transaction_session_timeout, lock_timeout, statement_timeout, track_activities, track_counts, and track_functions.

- Fix race condition that led to delayed delivery of interprocess signals on Windows (Amit Kapila)
  
  This caused visible timing oddities in NOTIFY, and perhaps other misbehavior.

- On Windows, retry a few times after an ERROR_ACCESS_DENIED file access failure (Alexander Lakhin, Tom Lane)
  
  This helps cope with cases where a file open attempt fails because the targeted file is flagged for deletion but not yet actually gone. pg_ctl, for example, frequently failed with such an error when probing to see if the postmaster had shut down yet.

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**E.2. Release 9.6.16**

**Release date:** 2019-11-14

This release contains a variety of fixes from 9.6.15. For information about new features in the 9.6 major release, see Section E.18.
E.2.1. Migration to Version 9.6.16

A dump/restore is not required for those running 9.6.X. However, if you use the contrib/intarray extension with a GiST index, and you rely on indexed searches for the $<$ operator, see the entry below about that. Also, if you are upgrading from a version earlier than 9.6.9, see Section E.9.

E.2.2. Changes

- Fix failure of `ALTER TABLE SET` with a custom relation option (Michael Paquier)
- Disallow changing a multiply-inherited column’s type if not all parent tables were changed (Tom Lane)
  
  Previously, this was allowed, whereupon queries on the now-out-of-sync parent would fail.
- Prevent `VACUUM` from trying to freeze an old multixact ID involving a still-running transaction (Nathan Bossart, Jeremy Schneider)
  
  This case would lead to `VACUUM` failing until the old transaction terminates.
- Ensure that offset expressions in `WINDOW` clauses are processed when a query’s expressions are manipulated (Andrew Gierth)
  
  This oversight could result in assorted failures when the offsets are nontrivial expressions. One example is that a function parameter reference in such an expression would fail if the function was inlined.
- Fix handling of whole-row variables in `WITH CHECK OPTION` expressions and row-level-security policy expressions (Andres Freund)
  
  Previously, such usage might result in bogus errors about row type mismatches.
- Avoid postmaster failure if a parallel query requests a background worker when no postmaster child process array slots remain free (Tom Lane)
- Prevent possible double-free if a `BEFORE UPDATE` trigger returns the old tuple as-is, and it is not the last such trigger (Thomas Munro)
- Provide a relevant error context line when an error occurs while setting GUC parameters during parallel worker startup (Thomas Munro)
- In serializable mode, ensure that row-level predicate locks are acquired on the correct version of the row (Thomas Munro, Heikki Linnakangas)
  
  If the visible version of the row is HOT-updated, the lock might be taken on its now-dead predecessor, resulting in subtle failures to guarantee serialization.
- Ensure that `fsync()` is applied only to files that are opened read/write (Andres Freund, Michael Paquier)
  
  Some code paths tried to do this after opening a file read-only, but on some platforms that causes “bad file descriptor” or similar errors.
- Allow encoding conversion to succeed on longer strings than before (Álvaro Herrera, Tom Lane)
  
  Previously, there was a hard limit of 0.25GB on the input string, but now it will work as long as the converted output is not over 1GB.
- Avoid creating unnecessarily-bulky tuple stores for window functions (Andrew Gierth)
Appendix E. Release Notes

In some cases the tuple storage would include all columns of the source table(s), not just the ones that are needed by the query.

- Allow `realloc()` to give back space when a large chunk is reduced in size (Tom Lane)
- Ensure that temporary WAL and history files are removed at the end of archive recovery (Sawada Masahiko)
- Avoid failure in archive recovery if `recovery_min_apply_delay` is enabled (Fujii Masao)
  - `recovery_min_apply_delay` is not typically used in this configuration, but it should work.
- Avoid unwanted delay during shutdown of a logical replication walsender (Craig Ringer, Álvaro Herrera)
- Correctly time-stamp replication messages for logical decoding (Jeff Janes)
  - This oversight resulted, for example, in `pg_stat_subscription.last_msg_send_time` usually reading as NULL.
- In logical decoding, ensure that sub-transactions are correctly accounted for when reconstructing a snapshot (Masahiko Sawada)
  - This error leads to assertion failures; it’s unclear whether any bad effects exist in production builds.
- Fix race condition during backend exit, when the backend process has previously waited for synchronous replication to occur (Dongming Liu)
- Fix `ALTER SYSTEM` to cope with duplicate entries in `postgresql.auto.conf` (Ian Barwick)
  - `ALTER SYSTEM` itself will not generate such a state, but external tools that modify `postgresql.auto.conf` could do so. Duplicate entries for the target variable will now be removed, and then the new setting (if any) will be appended at the end.
- Reject include directives with empty file names in configuration files, and report include-file recursion more clearly (Ian Barwick, Tom Lane)
- Avoid logging complaints about abandoned connections when using PAM authentication (Tom Lane)
  - libpq-based clients will typically make two connection attempts when a password is required, since they don’t prompt their user for a password until their first connection attempt fails. Therefore the server is coded not to generate useless log spam when a client closes the connection upon being asked for a password. However, the PAM authentication code hadn’t gotten that memo, and would generate several messages about a phantom authentication failure.
- Fix some cases where an incomplete date specification is not detected in `time with time zone` input (Alexander Lakhin)
  - If a time zone with a time-varying UTC offset is specified, then a date must be as well, so that the offset can be resolved. Depending on the syntax used, this check was not enforced in some cases, allowing bogus output to be produced.
- Fix misbehavior of `bitshiftright()` (Tom Lane)
  - The bitstring right shift operator failed to zero out padding space that exists in the last byte of the result when the bitstring length is not a multiple of 8. While invisible to most operations, any nonzero bits there would result in unexpected comparison behavior, since bitstring comparisons don’t bother to ignore the extra bits, expecting them to always be zero.
  - If you have inconsistent data as a result of saving the output of `bitshiftright()` in a table, it’s possible to fix it with something like
    
    ```sql
    UPDATE mytab SET bitcol = ~(~bitcol) WHERE bitcol != ~(~bitcol);
    ```
Appendix E. Release Notes

- Fix detection of edge-case integer overflow in interval multiplication (Yuya Watari)
- Avoid crashes if `ispell` text search dictionaries contain wrong affix data (Arthur Zakirov)
- Fix incorrect compression logic for GIN posting lists (Heikki Linnakangas)
  A GIN posting list item can require 7 bytes if the distance between adjacent indexed TIDs exceeds 16TB. One step in the logic was out of sync with that, and might try to write the value into a 6-byte buffer. In principle this could cause a stack overrun, but on most architectures it’s likely that the next byte would be unused alignment padding, making the bug harmless. In any case the bug would be very difficult to hit.
- Fix handling of infinity, NaN, and NULL values in KNN-GiST (Alexander Korotkov)
  The query’s output order could be wrong (different from a plain sort’s result) if some distances computed for non-null column values are infinity or NaN.
- Fix handling of searches for NULL in KNN-SP-GiST (Nikita Glukhov)
- On Windows, recognize additional spellings of the “Norwegian (Bokmål)” locale name (Tom Lane)
- Avoid compile failure if an ECPG client includes `ecpglib.h` while having `ENABLE_NLS` defined (Tom Lane)
  This risk was created by a misplaced declaration: `ecpg_gettext()` should not be visible to client code.
- In psql, resynchronize internal state about the server after an unexpected connection loss and successful reconnection (Peter Billen, Tom Lane)
  Ordinarily this is unnecessary since the state would be the same anyway. But it can matter in corner cases, such as where the connection might lead to one of several servers. This change causes psql to re-issue any interactive messages that it would have issued at startup, for example about whether SSL is in use.
- Avoid platform-specific null pointer dereference in psql (Quentin Rameau)
- Fix `pg_dump`’s handling of circular dependencies in views (Tom Lane)
  In some cases a view may depend on an object that `pg_dump` needs to dump later than the view; the most common example is that a query using `GROUP BY` on a primary-key column may be semantically invalid without the primary key. This is now handled by emitting a dummy `CREATE VIEW` command that just establishes the view’s column names and types, and then later emitting `CREATE OR REPLACE VIEW` with the full view definition. Previously, the dummy definition was actually a `CREATE TABLE` command, and this was automatically converted to a view by a later `CREATE RULE` command. The new approach has been used successfully in PostgreSQL version 10 and later. We are back-patching it into older releases now because of reports that the previous method causes bogus error messages about the view’s replica identity status. This change also avoids problems when trying to use the `--clean` option during a restore involving such a view.
- In `pg_dump`, ensure stable output order for similarly-named triggers and row-level-security policy objects (Benjie Gillam)
  Previously, if two triggers on different tables had the same names, they would be sorted in OID-based order, which is less desirable than sorting them by table name. Likewise for RLS policies.
- Fix `pg_dump` to work again with pre-8.3 source servers (Tom Lane)
  A previous fix caused `pg_dump` to always try to query `pg_opfamily`, but that catalog doesn’t exist before version 8.3.
- In `pg_restore`, treat `--f` as meaning “output to stdout” (Álvaro Herrera)
This synchronizes pg_restore’s behavior with some other applications, and in particular makes pre-v12 branches act similarly to version 12’s pg_restore, simplifying creation of dump/restore scripts that work across multiple PostgreSQL versions. Before this change, pg_restore interpreted such a switch as meaning “output to a file named ‘-‘”, but few people would want that.

• Improve pg_upgrade’s checks for the use of a data type that has changed representation, such as line (Tomas Vondra)
   The previous coding could be fooled by cases where the data type of interest underlies a stored column of a domain or composite type.

• Detect file read errors during pg_basebackup (Jeevan Chalke)

• In pg_rewind with an online source cluster, disable timeouts, much as pg_dump does (Alexander Kukushkin)

• Fix failure in pg_waldump with the -s option, when a continuation WAL record ends exactly at a page boundary (Andrey Lepikhov)

• In pg_waldump, include the newitemoff field in btree page split records (Peter Geoghegan)

• In pg_waldump with the --b kp-details option, avoid emitting extra newlines for WAL records involving full-page writes (Andres Freund)

• Fix small memory leak in pg_waldump (Andres Freund)

• Fix vacuumdb with a high --jobs option to handle running out of file descriptors better (Michael Paquier)

• Fix contrib/intarray’s GiST opclasses to not fail for empty arrays with <@ (Tom Lane)
   A clause like array_column <@ constant_array is considered indexable, but the index search may not find empty array values; of course, such entries should trivially match the search.
   The only practical back-patchable fix for this requires making <@ index searches scan the whole index, which is what this patch does. This is unfortunate: it means that the query performance is likely worse than a plain sequential scan would be.

   Applications whose performance is adversely impacted by this change have a couple of options. They could switch to a GIN index, which doesn’t have this bug, or they could replace array_column <@ constant_array with array_column <@ constant_array AND array_column && constant_array. That will provide about the same performance as before, and it will find all non-empty subsets of the given constant array, which is all that could reliably be expected of the query before.

• Allow configure --with-python to succeed when only python3 or only python2 can be found (Peter Eisenstraut, Tom Lane)
   Search for python, then python3, then python2, so that configure can succeed in the increasingly-more-common situation where there is no executable named simply python. It’s still possible to override this choice by setting the PYTHON environment variable.

• Fix configure’s test for presence of libperl so that it works on recent Red Hat releases (Tom Lane)
   Previously, it could fail if the user sets CFLAGS to -00.

• Ensure correct code generation for spinlocks on PowerPC (Noah Misch)
   The previous spinlock coding allowed the compiler to select register zero for use with an assembly instruction that does not accept that register, causing a build failure. We have seen only one long-ago report that matches this bug, but it could cause problems for people trying to build modified PostgreSQL code or use atypical compiler options.

• On PowerPC, avoid depending on the xlc compiler’s __fetch_and_add() function (Noah Misch)
Appendix E. Release Notes

xlc 13 and newer interpret this function in a way incompatible with our usage, resulting in an unusable build of PostgreSQL. Fix by using custom assembly code instead.

- On AIX, don’t use the compiler option -qsrcmsg (Noah Misch)
  This avoids an internal compiler error with xlc v16.1.0, with little consequence other than changing the format of compiler error messages.
- Fix MSVC build process to cope with spaces in the file path of OpenSSL (Andrew Dunstan)
- Update time zone data files to tzdata release 2019c for DST law changes in Fiji and Norfolk Island, plus historical corrections for Alberta, Austria, Belgium, British Columbia, Cambodia, Hong Kong, Indiana (Perry County), Kaliningrad, Kentucky, Michigan, Norfolk Island, South Korea, and Turkey.

E.3. Release 9.6.15

Release date: 2019-08-08

This release contains a variety of fixes from 9.6.14. For information about new features in the 9.6 major release, see Section E.18.

E.3.1. Migration to Version 9.6.15

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.9, see Section E.9.

E.3.2. Changes

- Require schema qualification to cast to a temporary type when using functional cast syntax (Noah Misch)
  We have long required invocations of temporary functions to explicitly specify the temporary schema, that is pg_temp.func_name(args). Require this as well for casting to temporary types using functional notation, for example pg_temp.type_name(arg). Otherwise it’s possible to capture a function call using a temporary object, allowing privilege escalation in much the same ways that we blocked in CVE-2007-2138. (CVE-2019-10208)
- Fix failure of ALTER TABLE ... ALTER COLUMN TYPE when altering multiple columns’ types in one command (Tom Lane)
  This fixes a regression introduced in the most recent minor releases: indexes using the altered columns were not processed correctly, leading to strange failures during ALTER TABLE.
- Don’t optimize away GROUP BY columns when the table involved is an inheritance parent (David Rowley)
  Normally, if a table’s primary key column(s) are included in GROUP BY, it’s safe to drop any other grouping columns, since the primary key columns are enough to make the groups unique. This rule does not work if the query is also reading inheritance child tables, though; the parent’s uniqueness does not extend to the children.
Appendix E. Release Notes

- Avoid using unnecessary sort steps for some queries with \texttt{GROUPING SETS} (Andrew Gierth, Richard Guo)
- Fix mishandling of multi-column foreign keys when rebuilding a foreign key constraint (Tom Lane) \texttt{ALTER TABLE} could make an incorrect decision about whether revalidation of a foreign key is necessary, if not all columns of the key are of the same type. It seems likely that the error would always have been in the conservative direction, that is revalidating unnecessarily.
- Avoid spurious deadlock errors when upgrading a tuple lock (Oleksii Kliukin)
  When two or more transactions are waiting for a transaction T1 to release a tuple-level lock, and T1 upgrades its lock to a higher level, a spurious deadlock among the waiting transactions could be reported when T1 finishes.
- Fix failure to resolve deadlocks involving multiple parallel worker processes (Rui Hai Jiang)
  It is not clear whether this bug is reachable with non-artificial queries, but if it did happen, the queries involved in an otherwise-resolvable deadlock would block until canceled.
- Prevent incorrect canonicalization of date ranges with \texttt{infinity} endpoints (Laurenz Albe)
  It’s incorrect to try to convert an open range to a closed one or vice versa by incrementing or decrementing the endpoint value, if the endpoint is infinite; so leave the range alone in such cases.
- Fix loss of fractional digits when converting very large \texttt{money} values to \texttt{numeric} (Tom Lane)
- Fix spinlock assembly code for MIPS CPUs so that it works on MIPS r6 (YunQiang Su)
- Make \texttt{libpq} ignore carriage return (\texttt{\r}) in connection service files (Tom Lane, Michael Paquier)
  In some corner cases, service files containing Windows-style newlines could be mis-parsed, resulting in connection failures.
- In \texttt{psql}, avoid offering incorrect tab completion options after \texttt{SET variable =} (Tom Lane)
- Fix \texttt{pg\_dump} to ensure that custom operator classes are dumped in the right order (Tom Lane)
  If a user-defined opclass is the subtype opclass of a user-defined range type, related objects were dumped in the wrong order, producing an unrestorable dump. (The underlying failure to handle opclass dependencies might manifest in other cases too, but this is the only known case.)
- Fix \texttt{contrib/passwordcheck} to coexist with other users of \texttt{check\_password\_hook} (Michael Paquier)
- Fix \texttt{contrib/sepgsql} tests to work under recent SELinux releases (Mike Palmiotto)
- Improve stability of \texttt{src/test/recovery} regression tests (Michael Paquier)
- Reduce stderr output from \texttt{pg\_upgrade}’s test script (Tom Lane)
- Fix TAP tests to work with msys Perl, in cases where the build directory is on a non-root msys mount point (Noah Misch)
- Support building Postgres with Microsoft Visual Studio 2019 (Haribabu Kommi)
- In Visual Studio builds, honor \texttt{WindowsSDKVersion} environment variable, if that’s set (Peifeng Qiu)
  This fixes build failures in some configurations.
- Support OpenSSL 1.1.0 and newer in Visual Studio builds (Juan José Santamaría Flecha, Michael Paquier)
- Allow make options to be passed down to \texttt{gmake} when non-GNU make is invoked at the top level (Thomas Munro)
- Avoid choosing \texttt{localtime} or \texttt{posixrules} as \texttt{TimeZone} during \texttt{initdb} (Tom Lane)
In some cases initdb would choose one of these artificial zone names over the “real” zone name. Prefer any other match to the C library’s timezone behavior over these two.

- Adjust `pg_timezone_names` view to show the Factory time zone if and only if it has a short abbreviation (Tom Lane)

Historically, IANA set up this artificial zone with an “abbreviation” like `Local time zone must be set--see zic manual page`. Modern versions of the tzdb database show `+00` instead, but some platforms alter the data to show one or another of the historical phrases. Show this zone only if it uses the modern abbreviation.

- Sync our copy of the timezone library with IANA tzcode release 2019b (Tom Lane)

  This adds support for zic’s new `-b slim` option to reduce the size of the installed zone files. We are not currently using that, but may enable it in future.

- Update time zone data files to tzdata release 2019b for DST law changes in Brazil, plus historical corrections for Hong Kong, Italy, and Palestine.

### E.4. Release 9.6.14

**Release date:** 2019-06-20

This release contains a variety of fixes from 9.6.13. For information about new features in the 9.6 major release, see Section E.18.

#### E.4.1. Migration to Version 9.6.14

A dump/restore is not required for those running 9.6.X. However, if you are upgrading from a version earlier than 9.6.9, see Section E.9.

#### E.4.2. Changes

- Fix failure of `ALTER TABLE ... ALTER COLUMN TYPE` when the table has a partial exclusion constraint (Tom Lane)

- Fix failure of `COMMENT` command for comments on domain constraints (Daniel Gustafsson, Michael Paquier)

- Fix faulty generation of merge-append plans (Tom Lane)

  This mistake could lead to “could not find pathkey item to sort” errors.

- Fix incorrect printing of queries with duplicate join names (Philip Dubé)

  This oversight caused a dump/restore failure for views containing such queries.

- Fix misoptimization of `{1,1}` quantifiers in regular expressions (Tom Lane)

  Such quantifiers were treated as no-ops and optimized away; but the documentation specifies that they impose greediness, or non-greediness in the case of the non-greedy variant `{1,1}?`, on the subexpression they’re attached to, and this did not happen. The misbehavior occurred only if the subexpression contained capturing parentheses or a back-reference.

- Avoid possible failures while initializing a new process’s `pg_stat_activity` data (Tom Lane)
Certain operations that could fail, such as converting strings extracted from an SSL certificate into the database encoding, were being performed inside a critical section. Failure there would result in database-wide lockup due to violating the access protocol for shared `pg_stat_activity` data.

- Fix race condition in check to see whether a pre-existing shared memory segment is still in use by a conflicting postmaster (Tom Lane)

- Avoid attempting to do database accesses for parameter checking in processes that are not connected to a specific database (Vignesh C, Andres Freund)

  This error could result in failures like “cannot read pg_class without having selected a database”.

- Avoid possible hang in `libpq` if using SSL and OpenSSL’s pending-data buffer contains an exact multiple of 256 bytes (David Binderman)

- Improve initdb’s handling of multiple equivalent names for the system time zone (Tom Lane, Andrew Gierth)

  Make initdb examine the `/etc/localtime` symbolic link, if that exists, to break ties between equivalent names for the system time zone. This makes initdb more likely to select the time zone name that the user would expect when multiple identical time zones exist. It will not change the behavior if `/etc/localtime` is not a symlink to a zone data file, nor if the time zone is determined from the `TZ` environment variable.

  Separately, prefer `UTC` over other spellings of that time zone, when neither `TZ` nor `/etc/localtime` provide a hint. This fixes an annoyance introduced by tzdata 2019a’s change to make the `UCT` and `UTC` zone names equivalent: initdb was then preferring `UCT`, which almost nobody wants.

- Fix ordering of `GRANT` commands emitted by `pg_dump` and `pg_dumpall` for databases and tablespaces (Nathan Bossart, Michael Paquier)

  If cascading grants had been issued, restore might fail due to the `GRANT` commands being given in an order that didn’t respect their interdependencies.

- Fix misleading error reports from `reindexdb` (Julien Rouhaud)

- Ensure that `vaccumdb` returns correct status if an error occurs while using parallel jobs (Julien Rouhaud)

- Fix `contrib/auto_explain` to not cause problems in parallel queries (Tom Lane)

  Previously, a parallel worker might try to log its query even if the parent query were not being logged by `auto_explain`. This would work sometimes, but it’s confusing, and in some cases it resulted in failures like “could not find key N in shm TOC”.

  Also, fix an off-by-one error that resulted in not necessarily logging every query even when the sampling rate is set to 1.0.

- In `contrib/postgres_fdw`, account for possible data modifications by local `BEFORE ROW UPDATE` triggers (Shohei Mochizuki)

  If a trigger modified a column that was otherwise not changed by the `UPDATE`, the new value was not transmitted to the remote server.

- On Windows, avoid failure when the database encoding is set to `SQL_ASCII` and we attempt to log a non-ASCII string (Noah Misch)

  The code had been assuming that such strings must be in UTF-8, and would throw an error if they didn’t appear to be validly encoded. Now, just transmit the untranslated bytes to the log.

- Make PL/pgSQL’s header files C++-safe (George Tarasov)
E.5. Release 9.6.13

Release date: 2019-05-09

This release contains a variety of fixes from 9.6.12. For information about new features in the 9.6 major release, see Section E.18.

E.5.1. Migration to Version 9.6.13

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.9, see Section E.9.

E.5.2. Changes

- Prevent row-level security policies from being bypassed via selectivity estimators (Dean Rasheed)

  Some of the planner’s selectivity estimators apply user-defined operators to values found in pg_statistic (e.g., most-common values). A leaky operator therefore can disclose some of the entries in a data column, even if the calling user lacks permission to read that column. In CVE-2017-7484 we added restrictions to forestall that, but we failed to consider the effects of row-level security. A user who has SQL permission to read a column, but who is forbidden to see certain rows due to RLS policy, might still learn something about those rows’ contents via a leaky operator. This patch further tightens the rules, allowing leaky operators to be applied to statistics data only when there is no relevant RLS policy. (CVE-2019-10130)

- Fix behavior for an UPDATE or DELETE on an inheritance tree or partitioned table in which every table can be excluded (Amit Langote, Tom Lane)

  In such cases, the query did not report the correct set of output columns when a RETURNING clause was present, and if there were any statement-level triggers that should be fired, it didn’t fire them.

- Fix handling of explicit DEFAULT items in an INSERT ... VALUES command with multiple VALUES rows, if the target relation is an updatable view (Amit Langote, Dean Rasheed)

  When the updatable view has no default for the column but its underlying table has one, a single-row INSERT ... VALUES will use the underlying table’s default. In the multi-row case, however, NULL was always used. Correct it to act like the single-row case.

- Fix CREATE VIEW to allow zero-column views (Ashutosh Sharma)

  We should allow this for consistency with allowing zero-column tables. Since a table can be converted to a view, zero-column views could be created even with the restriction in place, leading to dump/reload failures.

- Add missing support for CREATE TABLE IF NOT EXISTS ... AS EXECUTE ... (Andreas Karlsson)

  The combination of IF NOT EXISTS and EXECUTE should work, but the grammar omitted it.

- Ensure that sub-SELECTs appearing in row-level-security policy expressions are executed with the correct user’s permissions (Dean Rasheed)

  Previously, if the table having the RLS policy was accessed via a view, such checks might be executed as the user calling the view, not as the view owner as they should be.

- Accept XML documents as valid values of type xml when xmloption is set to content, as required by SQL:2006 and later (Chapman Flack)
Appendix E. Release Notes

Previously PostgreSQL followed the SQL:2003 definition, which doesn’t allow this. But that creates a serious problem for dump/restore: there is no setting of `xmloption` that will accept all valid XML data. Hence, switch to the 2006 definition.

```
pg_dump is also modified to emit `SET xmloption = content` while restoring data, ensuring that dump/restore works even if the prevailing setting is `document`.
```

- Improve server’s startup-time checks for whether a pre-existing shared memory segment is still in use (Noah Misch)

  The postmaster is now more likely to detect that there are still active processes from a previous postmaster incarnation, even if the `postmaster.pid` file has been removed.

- Avoid counting parallel workers’ transactions as separate transactions (Haribabu Kommi)

- Fix incompatibility of GIN-index WAL records (Alexander Korotkov)

  A fix applied in February’s minor releases was not sufficiently careful about backwards compatibility, leading to problems if a standby server of that vintage reads GIN page-deletion WAL records generated by a primary server of a previous minor release.

- Tolerate `EINVAL` and `ENOSYS` error results, where appropriate, for `fsync` and `sync_file_range` calls (Thomas Munro, James Sewell)

  The previous change to panic on file synchronization failures turns out to have been excessively paranoid for certain cases where a failure is predictable and essentially means “operation not supported”.

- Fix “failed to build any N-way joins” planner failures with lateral references leading out of FULL outer joins (Tom Lane)

- Check the appropriate user’s permissions when enforcing rules about letting a leaky operator see `pg_statistic` data (Dean Rasheed)

  When an underlying table is being accessed via a view, consider the privileges of the view owner while deciding whether leaky operators may be applied to the table’s statistics data, rather than the privileges of the user making the query. This makes the planner’s rules about what data is visible match up with the executor’s, avoiding unnecessarily-poor plans.

- Speed up planning when there are many equality conditions and many potentially-relevant foreign key constraints (David Rowley)

- Avoid O(N^2) performance issue when rolling back a transaction that created many tables (Tomas Vondra)

- Fix race conditions in management of dynamic shared memory (Thomas Munro)

  These could lead to “dsa_area could not attach to segment” or “cannot unpin a segment that is not pinned” errors.

- Fix race condition in which a hot-standby postmaster could fail to shut down after receiving a smart-shutdown request (Tom Lane)

- Fix possible crash when `pg_identify_object_as_address()` is given invalid input (Álvaro Herrera)

- Tighten validation of encoded SCRAM-SHA-256 and MD5 passwords (Jonathan Katz)

  A password string that had the right initial characters could be mistaken for one that is correctly hashed into SCRAM-SHA-256 or MD5 format. The password would be accepted but would be unusable later.

- Fix handling of `lc_time` settings that imply an encoding different from the database’s encoding (Juan José Santamaría Flecha, Tom Lane)
Localized month or day names that include non-ASCII characters previously caused unexpected errors or wrong output in such locales.

- Fix incorrect `operator_precedence_warning` checks involving unary minus operators (Rikard Falkeborn)
- Disallow `NaN` as a value for floating-point server parameters (Tom Lane)
- Rearrange `REINDEX` processing to avoid assertion failures when reindexing individual indexes of `pg_class` (Andres Freund, Tom Lane)
- Fix planner assertion failure for parameterized dummy paths (Tom Lane)
- Insert correct test function in the result of `SnapBuildInitialSnapshot()` (Antonin Houska)
  No core code cares about this, but some extensions do.
- Fix intermittent “could not reattach to shared memory” session startup failures on Windows (Noah Misch)
  A previously unrecognized source of these failures is creation of thread stacks for a process’s default thread pool. Arrange for such stacks to be allocated in a different memory region.
- Fix error detection in directory scanning on Windows (Konstantin Knizhnik)
  Errors, such as lack of permissions to read the directory, were not detected or reported correctly; instead the code silently acted as though the directory were empty.
- Fix grammar problems in ecpg (Tom Lane)
  A missing semicolon led to mistranslation of `SET variable = DEFAULT` (but not `SET variable TO DEFAULT`) in ecpg programs, producing syntactically invalid output that the server would reject.
  Additionally, in a `DROP TYPE` or `DROP DOMAIN` command that listed multiple type names, only the first type name was actually processed.
- Sync ecpg’s syntax for `CREATE TABLE AS` with the server’s (Daisuke Higuchi)
- Fix possible buffer overruns in ecpg’s processing of include filenames (Liu Huailing, Fei Wu)
- Avoid crash in `contrib/vacuumlo` if an `lo_unlink()` call failed (Tom Lane)
- Sync our copy of the timezone library with IANA tzcode release 2019a (Tom Lane)
  This corrects a small bug in zic that caused it to output an incorrect year-2440 transition in the `Africa/Casablanca` zone, and adds support for zic’s new `-r` option.
- Update time zone data files to tzdata release 2019a for DST law changes in Palestine and Metlakatla, plus historical corrections for Israel.
  `Etc/UCT` is now a backward-compatibility link to `Etc/UTC`, instead of being a separate zone that generates the abbreviation `UCT`, which nowadays is typically a typo. PostgreSQL will still accept `UCT` as an input zone abbreviation, but it won’t output it.

**E.6. Release 9.6.12**

**Release date:** 2019-02-14

This release contains a variety of fixes from 9.6.11. For information about new features in the 9.6 major release, see Section E.18.
Appendix E. Release Notes


A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.9, see Section E.9.

E.6.2. Changes

- By default, panic instead of retrying after \texttt{fsync()} failure, to avoid possible data corruption (Craig Ringer, Thomas Munro)

Some popular operating systems discard kernel data buffers when unable to write them out, reporting this as \texttt{fsync()} failure. If we reissue the \texttt{fsync()} request it will succeed, but in fact the data has been lost, so continuing risks database corruption. By raising a panic condition instead, we can replay from WAL, which may contain the only remaining copy of the data in such a situation. While this is surely ugly and inefficient, there are few alternatives, and fortunately the case happens very rarely.

A new server parameter \texttt{data_sync_retry} has been added to control this; if you are certain that your kernel does not discard dirty data buffers in such scenarios, you can set \texttt{data_sync_retry} to \texttt{on} to restore the old behavior.

- Include each major release branch’s release notes in the documentation for only that branch, rather than that branch and all later ones (Tom Lane)

The duplication induced by the previous policy was getting out of hand. Our plan is to provide a full archive of release notes on the project’s web site, but not duplicate it within each release.

- Avoid possible deadlock when acquiring multiple buffer locks (Nishant Fnu)

- Avoid deadlock between hot-standby queries and replay of GIN index page deletion (Alexander Korotkov)

- Fix possible crashes in logical replication when index expressions or predicates are in use (Peter Eisentraut)

- Avoid useless and expensive logical decoding of TOAST data during a table rewrite (Tomas Vondra)

- Fix logic for stopping a subset of WAL senders when synchronous replication is enabled (Paul Guo, Michael Paquier)

- Avoid possibly writing an incorrect replica identity field in a tuple deletion WAL record (Stas Kelvich)

- Make the archiver prioritize WAL history files over WAL data files while choosing which file to archive next (David Steele)

- Fix possible crash in \texttt{UPDATE} with a multiple \texttt{SET} clause using a sub-\texttt{SELECT} as source (Tom Lane)

- Avoid crash if libxml2 returns a null error message (Sergio Conde Gómez)

- Fix spurious grouping-related parser errors caused by inconsistent handling of collation assignment (Andrew Gierth)

In some cases, expressions that should be considered to match were not seen as matching, if they included operations on collatable data types.

- Check whether the comparison function underlying \texttt{LEAST()} or \texttt{GREATEST()} is leakproof, rather than just assuming it is (Tom Lane)
Actual information leaks from btree comparison functions are typically hard to provoke, but in principle they could happen.

- **Fix incorrect planning of queries involving nested loops both above and below a Gather plan node (Tom Lane)**
  If both levels of nestloop needed to pass the same variable into their right-hand sides, an incorrect plan would be generated.

- **Fix incorrect planning of queries in which a lateral reference must be evaluated at a foreign table scan (Tom Lane)**

- **Fix corner-case underestimation of the cost of a merge join (Tom Lane)**
  The planner could prefer a merge join when the outer key range is much smaller than the inner key range, even if there are so many duplicate keys on the inner side that this is a poor choice.

- **Avoid O(N^2) planning time growth when a query contains many thousand indexable clauses (Tom Lane)**

- **Improve `ANALYZE`’s handling of concurrently-updated rows (Jeff Janes, Tom Lane)**
  Previously, rows deleted by an in-progress transaction were omitted from `ANALYZE`’s sample, but this has been found to lead to more inconsistency than including them would do. In effect, the sample now corresponds to an MVCC snapshot as of `ANALYZE`’s start time.

- **Make `TRUNCATE` ignore inheritance child tables that are temporary tables of other sessions (Amit Langote, Michael Paquier)**
  This brings `TRUNCATE` into line with the behavior of other commands. Previously, such cases usually ended in failure.

- **Fix `TRUNCATE` to update the statistics counters for the right table (Tom Lane)**
  If the truncated table had a TOAST table, that table’s counters were reset instead.

- **Process `ALTER TABLE ONLY ADD COLUMN IF NOT EXISTS` correctly (Greg Stark)**

- **Allow `UNLISTEN` in hot-standby mode (Shay Rojansky)**
  This is necessarily a no-op, because `LISTEN` isn’t allowed in hot-standby mode; but allowing the dummy operation simplifies session-state-reset logic in clients.

- **Fix missing role dependencies in some schema and data type permissions lists (Tom Lane)**
  In some cases it was possible to drop a role to which permissions had been granted. This caused no immediate problem, but a subsequent dump/reload or upgrade would fail, with symptoms involving attempts to grant privileges to all-numeric role names.

- **Ensure relation caches are updated properly after adding or removing foreign key constraints (Álvaro Herrera)**
  This oversight could result in existing sessions failing to enforce a newly-created constraint, or continuing to enforce a dropped one.

- **Ensure relation caches are updated properly after renaming constraints (Amit Langote)**

- **Make autovacuum more aggressive about removing leftover temporary tables, and also remove leftover temporary tables during `DISCARD TEMP` (Álvaro Herrera)**
  This helps ensure that remnants from a crashed session are cleaned up more promptly.

- **Fix replay of GiST index micro-vacuum operations so that concurrent hot-standby queries do not see inconsistent state (Alexander Korotkov)**
- Prevent empty GIN index pages from being reclaimed too quickly, causing failures of concurrent searches (Andrey Borodin, Alexander Korotkov)

- Fix edge-case failures in float-to-integer coercions (Andrew Gierth, Tom Lane)
  Values very slightly above the maximum valid integer value might not be rejected, and then would overflow, producing the minimum valid integer instead. Also, values that should round to the minimum or maximum integer value might be incorrectly rejected.

- When making a PAM authentication request, don’t set the PAM_RHOST variable if the connection is via a Unix socket (Thomas Munro)
  Previously that variable would be set to [local], which is at best unhelpful, since it’s supposed to be a host name.

- Disallow setting client_min_messages higher than ERROR (Jonah Harris, Tom Lane)
  Previously, it was possible to set this variable to FATAL or PANIC, which had the effect of suppressing transmission of ordinary error messages to the client. However, that’s contrary to guarantees that are given in the PostgreSQL wire protocol specification, and it caused some clients to become very confused. In released branches, fix this by silently treating such settings as meaning ERROR instead. Version 12 and later will reject those alternatives altogether.

- Fix ecpglib to use uselocale() or _configthreadlocale() in preference to setlocale() (Michael Meskes, Tom Lane)
  Since setlocale() is not thread-local, and might not even be thread-safe, the previous coding caused problems in multi-threaded ecpg applications.

- Fix incorrect results for numeric data passed through an ecpg SQLDA (SQL Descriptor Area) (Daisuke Higuchi)
  Values with leading zeroes were not copied correctly.

- Fix psql’s \g target meta-command to work with COPY TO STDOUT (Daniel Vérité)
  Previously, the target option was ignored, so that the copy data always went to the current query output target.

- Make psql’s LaTeX output formats render special characters properly (Tom Lane)
  Backslash and some other ASCII punctuation characters were not rendered correctly, leading to document syntax errors or wrong characters in the output.

- Fix pg_dump’s handling of materialized views with indirect dependencies on primary keys (Tom Lane)
  This led to mis-labeling of such views’ dump archive entries, causing harmless warnings about “archive items not in correct section order”; less harmlessly, selective-restore options depending on those labels, such as --section, might misbehave.

- Avoid null-pointer-dereference crash on some platforms when pg_dump or pg_restore tries to report an error (Tom Lane)

- Fix contrib/hstore to calculate correct hash values for empty hstore values that were created in version 8.4 or before (Andrew Gierth)
  The previous coding did not give the same result as for an empty hstore value created by a newer version, thus potentially causing wrong results in hash joins or hash aggregation. It is advisable to reindex any hash indexes built on hstore columns; if the table might contain data that was originally stored as far back as 8.4 and was never dumped/reloaded since then.

- Avoid crashes and excessive runtime with large inputs to contrib/intarray’s gist__int_ops index support (Andrew Gierth)
• Support new Makefile variables PG_CFLAGS, PG_CXXFLAGS, and PG_LDFLAGS in pgxs builds (Christoph Berg)
  This simplifies customization of extension build processes.
• Fix Perl-coded build scripts to not assume “.” is in the search path, since recent Perl versions don’t include that (Andrew Dunstan)
• Fix server command-line option parsing problems on OpenBSD (Tom Lane)
• Relocate call of set_rel_pathlist_hook so that extensions can use it to supply partial paths for parallel queries (KaiGai Kohei)
  This is not expected to affect existing use-cases.
• Update time zone data files to tzdata release 2018i for DST law changes in Kazakhstan, Metlakatla, and Sao Tome and Principe. Kazakhstan’s Qyzylorda zone is split in two, creating a new zone Asia/Qostanay, as some areas did not change UTC offset. Historical corrections for Hong Kong and numerous Pacific islands.

E.7. Release 9.6.11

Release date: 2018-11-08
This release contains a variety of fixes from 9.6.10. For information about new features in the 9.6 major release, see Section E.18.

E.7.1. Migration to Version 9.6.11
A dump/restore is not required for those running 9.6.X.
However, if you are upgrading from a version earlier than 9.6.9, see Section E.9.

E.7.2. Changes
• Fix corner-case failures in has_foo_privilege() family of functions (Tom Lane)
  Return NULL rather than throwing an error when an invalid object OID is provided. Some of these functions got that right already, but not all. has_column_privilege() was additionally capable of crashing on some platforms.
• Avoid O(N^2) slowdown in regular expression match/split functions on long strings (Andrew Gierth)
• Fix parsing of standard multi-character operators that are immediately followed by a comment or + or − (Andrew Gierth)
  This oversight could lead to parse errors, or to incorrect assignment of precedence.
• Avoid O(N^3) slowdown in lexer for long strings of + or − characters (Andrew Gierth)
• Fix mis-execution of SubPlans when the outer query is being scanned backwards (Andrew Gierth)
• Fix failure of UPDATE/DELETE ... WHERE CURRENT OF ... after rewinding the referenced cursor (Tom Lane)
A cursor that scans multiple relations (particularly an inheritance tree) could produce wrong behavior if rewound to an earlier relation.

- **Fix EvalPlanQual to handle conditionally-executed InitPlans properly** (Andrew Gierth, Tom Lane)
  
  This resulted in hard-to-reproduce crashes or wrong answers in concurrent updates, if they contained code such as an uncorrelated sub-SELECT inside a CASE construct.

- **Fix character-class checks to not fail on Windows for Unicode characters above U+FFFF** (Tom Lane, Kenji Uno)
  
  This bug affected full-text-search operations, as well as contrib/ltree and contrib/pg_trgm.

- **Disallow pushing sub-SELECTs containing window functions, LIMIT, or OFFSET to parallel workers** (Amit Kapila)
  
  Such cases could result in inconsistent behavior due to different workers getting different answers, as a result of indeterminacy due to row-ordering variations.

- **Ensure that sequences owned by a foreign table are processed by ALTER OWNER on the table** (Peter Eisentraut)
  
  The ownership change should propagate to such sequences as well, but this was missed for foreign tables.

- **Ensure that the server will process already-received NOTIFY and SIGTERM interrupts before waiting for client input** (Jeff Janes, Tom Lane)

- **Fix over-allocation of space for array_out()’s result string** (Keiichi Hirobe)

- **Fix memory leak in repeated SP-GiST index scans** (Tom Lane)
  
  This is only known to amount to anything significant in cases where an exclusion constraint using SP-GiST receives many new index entries in a single command.

- **Ensure that ApplyLogicalMappingFile() closes the mapping file when done with it** (Tomas Vondra)
  
  Previously, the file descriptor was leaked, eventually resulting in failures during logical decoding.

- **Fix logical decoding to handle cases where a mapped catalog table is repeatedly rewritten, e.g. by VACUUM FULL** (Andres Freund)

- **Prevent starting the server with wal_level set to too low a value to support an existing replication slot** (Andres Freund)

- **Avoid crash if a utility command causes infinite recursion** (Tom Lane)

- **When initializing a hot standby, cope with duplicate XIDs caused by two-phase transactions on the master** (Michael Paquier, Konstantin Knizhnik)

- **Fix event triggers to handle nested ALTER TABLE commands** (Michael Paquier, Álvaro Herrera)

- **Propagate parent process’s transaction and statement start timestamps to parallel workers** (Konstantin Knizhnik)
  
  This prevents misbehavior of functions such as transaction_timestamp() when executed in a worker.

- **Fix transfer of expanded datums to parallel workers so that alignment is preserved, preventing crashes on alignment-picky platforms** (Tom Lane, Amit Kapila)

- **Fix WAL file recycling logic to work correctly on standby servers** (Michael Paquier)
Appendix E. Release Notes

Depending on the setting of `archive_mode`, a standby might fail to remove some WAL files that could be removed.

- Fix handling of commit-timestamp tracking during recovery (Masahiko Sawada, Michael Paquier)
  If commit timestamp tracking has been turned on or off, recovery might fail due to trying to fetch the commit timestamp for a transaction that did not record it.

- Randomize the `random()` seed in bootstrap and standalone backends, and in initdb (Noah Misch)
  The main practical effect of this change is that it avoids a scenario where initdb might mistakenly conclude that POSIX shared memory is not available, due to name collisions caused by always using the same random seed.

- Allow DSM allocation to be interrupted (Chris Travers)

- Avoid failure in a parallel worker when loading an extension that tries to access system caches within its init function (Thomas Munro)
  We don’t consider that to be good extension coding practice, but it mostly worked before parallel query, so continue to support it for now.

- Properly handle turning `full_page_writes` on dynamically (Kyotaro Horiguchi)

- Fix possible crash due to double `free()` during SP-GiST rescan (Andrew Gierth)

- Avoid possible buffer overrun when replaying GIN page recompression from WAL (Alexander Korotkov, Sivasubramanian Ramasubramanian)

- Fix missed `fsync` of a replication slot’s directory (Konstantin Knizhnik, Michael Paquier)

- Fix unexpected timeouts when using `wal_sender_timeout` on a slow server (Noah Misch)

- Ensure that hot standby processes use the correct WAL consistency point (Alexander Kukushkin, Michael Paquier)
  This prevents possible misbehavior just after a standby server has reached a consistent database state during WAL replay.

- Ensure background workers are stopped properly when the postmaster receives a fast-shutdown request before completing database startup (Alexander Kukushkin)

- Update the free space map during WAL replay of page all-visible/frozen flag changes (Álvaro Herrera)
  Previously we were not careful about this, reasoning that the FSM is not critical data anyway. However, if it’s sufficiently out of date, that can result in significant performance degradation after a standby has been promoted to primary. The FSM will eventually be healed by updates, but we’d like it to be good sooner, so work harder at maintaining it during WAL replay.

- Avoid premature release of parallel-query resources when query end or tuple count limit is reached (Amit Kapila)
  It’s only okay to shut down the executor at this point if the caller cannot demand backwards scan afterwards.

- Don’t run atexit callbacks when servicing `SIGQUIT` (Heikki Linnakangas)

- Don’t record foreign-server user mappings as members of extensions (Tom Lane)
  If `CREATE USER MAPPING` is executed in an extension script, an extension dependency was created for the user mapping, which is unexpected. Roles can’t be extension members, so user mappings shouldn’t be either.

- Make syslogger more robust against failures in opening CSV log files (Tom Lane)
• Fix psql, as well as documentation examples, to call PQconsumeInput() before each PQnotifies() call (Tom Lane)
  This fixes cases in which psql would not report receipt of a NOTIFY message until after the next command.
• Fix possible inconsistency in pg_dump’s sorting of dissimilar object names (Jacob Champion)
• Ensure that pg_restore will schema-qualify the table name when emitting DISABLE/ENABLE TRIGGER commands (Tom Lane)
  This avoids failures due to the new policy of running restores with restrictive search path.
• Fix pg_upgrade to handle event triggers in extensions correctly (Haribabu Kommi)
  pg_upgrade failed to preserve an event trigger’s extension-membership status.
• Fix pg_upgrade’s cluster state check to work correctly on a standby server (Bruce Momjian)
• Enforce type cube’s dimension limit in all contrib/cube functions (Andrey Borodin)
  Previously, some cube-related functions could construct values that would be rejected by cube_in(), leading to dump/reload failures.
• In contrib/postgres_fdw, don’t try to ship a variable-free ORDER BY clause to the remote server (Andrew Gierth)
• Fix contrib/unaccent’s unaccent() function to use the unaccent text search dictionary that is in the same schema as the function (Tom Lane)
  Previously it tried to look up the dictionary using the search path, which could fail if the search path has a restrictive value.
• Fix build problems on macOS 10.14 (Mojave) (Tom Lane)
  Adjust configure to add an -isysroot switch to CPPFLAGS; without this, PL/Perl and PL/Tcl fail to configure or build on macOS 10.14. The specific sysroot used can be overridden at configure time or build time by setting the PG_SYSROOT variable in the arguments of configure or make.
  It is now recommended that Perl-related extensions write $(perl_includespec) rather than -I$(perl_archlibexp)/CORE in their compiler flags. The latter continues to work on most platforms, but not recent macOS.
  Also, it should no longer be necessary to specify --with-tclconfig manually to get PL/Tcl to build on recent macOS releases.
• Fix MSVC build and regression-test scripts to work on recent Perl versions (Andrew Dunstan)
  Perl no longer includes the current directory in its search path by default; work around that.
• On Windows, allow the regression tests to be run by an Administrator account (Andrew Dunstan)
  To do this safely, pg_regress now gives up any such privileges at startup.
• Allow btree comparison functions to return INT_MIN (Tom Lane)
  Up to now, we’ve forbidden datatype-specific comparison functions from returning INT_MIN, which allows callers to invert the sort order just by negating the comparison result. However, this was never safe for comparison functions that directly return the result of memcmp(), strcmp(), etc, as POSIX doesn’t place any such restriction on those functions. At least some recent versions of memcmp() can return INT_MIN, causing incorrect sort ordering. Hence, we’ve removed this restriction. Callers must now use the INVERT_COMPARE_RESULT() macro if they wish to invert the sort order.
• Fix recursion hazard in shared-invalidation message processing (Tom Lane)
This error could, for example, result in failure to access a system catalog or index that had just been processed by VACUUM FULL.

This change adds a new result code for LockAcquire, which might possibly affect external callers of that function, though only very unusual usage patterns would have an issue with it. The API of LockAcquireExtended is also changed.

- Save and restore SPI’s global variables during SPI_connect() and SPI_finish() (Chapman Flack, Tom Lane)

  This prevents possible interference when one SPI-using function calls another.

- Avoid using potentially-under-aligned page buffers (Tom Lane)

  Invent new union types PGAlignedBlock and PGAlignedXLogBlock, and use these in place of plain char arrays, ensuring that the compiler can’t place the buffer at a misaligned start address. This fixes potential core dumps on alignment-picky platforms, and may improve performance even on platforms that allow misalignment.

- Make src/port/snprintf.c follow the C99 standard’s definition of snprintf()’s result value (Tom Lane)

  On platforms where this code is used (mostly Windows), its pre-C99 behavior could lead to failure to detect buffer overrun, if the calling code assumed C99 semantics.

- When building on i386 with the clang compiler, require -msse2 to be used (Andres Freund)

  This avoids problems with missed floating point overflow checks.

- Fix configure’s detection of the result type of strerror_r() (Tom Lane)

  The previous coding got the wrong answer when building with icc on Linux (and perhaps in other cases), leading to libpq not returning useful error messages for system-reported errors.

- Update time zone data files to tzdata release 2018g for DST law changes in Chile, Fiji, Morocco, and Russia (Volgograd), plus historical corrections for China, Hawaii, Japan, Macau, and North Korea.

E.8. Release 9.6.10

Release date: 2018-08-09

This release contains a variety of fixes from 9.6.9. For information about new features in the 9.6 major release, see Section E.18.

E.8.1. Migration to Version 9.6.10

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.9, see Section E.9.

E.8.2. Changes

- Fix failure to reset libpq’s state fully between connection attempts (Tom Lane)
An unprivileged user of dblink or postgres_fdw could bypass the checks intended to prevent use of server-side credentials, such as a ~/.pgpass file owned by the operating-system user running the server. Servers allowing peer authentication on local connections are particularly vulnerable. Other attacks such as SQL injection into a postgres_fdw session are also possible. Attacking postgres_fdw in this way requires the ability to create a foreign server object with selected connection parameters, but any user with access to dblink could exploit the problem. In general, an attacker with the ability to select the connection parameters for a libpq-using application could cause mischief, though other plausible attack scenarios are harder to think of. Our thanks to Andrew Krasichkov for reporting this issue. (CVE-2018-10915)

- **Fix** INSERT ... ON CONFLICT UPDATE through a view that isn’t just SELECT * FROM ... (Dean Rasheed, Amit Langote)

  Erroneous expansion of an updatable view could lead to crashes or “attribute ... has the wrong type” errors, if the view’s SELECT list doesn’t match one-to-one with the underlying table’s columns. Furthermore, this bug could be leveraged to allow updates of columns that an attacking user lacks UPDATE privilege for, if that user has INSERT and UPDATE privileges for some other column(s) of the table. Any user could also use it for disclosure of server memory. (CVE-2018-10925)

- **Ensure that updates to the relfrozenxid and relminmxid values for “nailed” system catalogs are processed in a timely fashion** (Andres Freund)

  Overoptimistic caching rules could prevent these updates from being seen by other sessions, leading to spurious errors and/or data corruption. The problem was significantly worse for shared catalogs, such as pg_authid, because the stale cache data could persist into new sessions as well as existing ones.

- **Fix case where a freshly-promoted standby crashes before having completed its first post-recovery checkpoint** (Michael Paquier, Kyotaro Horiguchi, Pavan Deolasee, Álvaro Herrera)

  This led to a situation where the server did not think it had reached a consistent database state during subsequent WAL replay, preventing restart.

- **Avoid emitting a bogus WAL record when recycling an all-zero btree page** (Amit Kapila)

  This mistake has been seen to cause assertion failures, and potentially it could result in unnecessary query cancellations on hot standby servers.

- **During WAL replay, guard against corrupted record lengths exceeding 1GB** (Michael Paquier)

  Treat such a case as corrupt data. Previously, the code would try to allocate space and get a hard error, making recovery impossible.

- **When ending recovery, delay writing the timeline history file as long as possible** (Heikki Linnakan-gas)

  This avoids some situations where a failure during recovery cleanup (such as a problem with a two-phase state file) led to inconsistent timeline state on-disk.

- **Improve performance of WAL replay for transactions that drop many relations** (Fujii Masao)

  This change reduces the number of times that shared buffers are scanned, so that it is of most benefit when that setting is large.

- **Improve performance of lock releasing in standby server WAL replay** (Thomas Munro)

- **Make logical WAL senders report streaming state correctly** (Simon Riggs, Sawada Masahiko)

  The code previously mis-detected whether or not it had caught up with the upstream server.

- **Fix bugs in snapshot handling during logical decoding, allowing wrong decoding results in rare cases** (Arseny Sher, Álvaro Herrera)
Appendix E. Release Notes

- Ensure a table’s cached index list is correctly rebuilt after an index creation fails partway through (Peter Geoghegan)

  Previously, the failed index’s OID could remain in the list, causing problems later in the same session.

- Fix mishandling of empty uncompressed posting list pages in GIN indexes (Sivasubramanian Rama-
  subramanian, Alexander Korotkov)

  This could result in an assertion failure after pg_upgrade of a pre-9.4 GIN index (9.4 and later will
  not create such pages).

- Ensure that VACUUM will respond to signals within btree page deletion loops (Andres Freund)

  Corrupted btree indexes could result in an infinite loop here, and that previously wasn’t interruptible
  without forcing a crash.

- Fix misoptimization of equivalence classes involving composite-type columns (Tom Lane)

  This resulted in failure to recognize that an index on a composite column could provide the sort
  order needed for a mergejoin on that column.

- Fix planner to avoid “ORDER/GROUP BY expression not found in targetlist” errors in some
queries with set-returning functions (Tom Lane)

- Fix SQL-standard FETCH FIRST syntax to allow parameters (\$n), as the standard expects (Andrew
  Gierth)

- Fix EXPLAIN’s accounting for resource usage, particularly buffer accesses, in parallel workers
(Amit Kapila, Robert Haas)

- Fix failure to schema-qualify some object names in getObjectDescription output (Kyotaro
Horiguchi, Tom Lane)

  Names of collations, conversions, and text search objects were not schema-qualified when they
should be.

- Fix CREATE AGGREGATE type checking so that parallelism support functions can be attached to
  variadic aggregates (Alexey Bashtanov)

- Widen COPY FROM’s current-line-number counter from 32 to 64 bits (David Rowley)

  This avoids two problems with input exceeding 4G lines: COPY FROM WITH HEADER would drop
  a line every 4G lines, not only the first line, and error reports could show a wrong line number.

- Add a string freeing function to ecpg’s pgtypes library, so that cross-module memory management
problems can be avoided on Windows (Takayuki Tsunakawa)

  On Windows, crashes can ensue if the free call for a given chunk of memory is not made from
the same DLL that malloc’ed the memory. The pgtypes library sometimes returns strings that it
expects the caller to free, making it impossible to follow this rule. Add a PGTYPESchar_free() function
that just wraps free, allowing applications to follow this rule.

- Fix ecpg’s support for long long variables on Windows, as well as other platforms that declare
strtol/strtoull nonstandardly or not at all (Dang Minh Huong, Tom Lane)

- Fix misidentification of SQL statement type in PL/pgSQL, when a rule change causes a change in
the semantics of a statement intra-session (Tom Lane)

  This error led to assertion failures, or in rare cases, failure to enforce the INTO STRICT option as
expected.

- Fix password prompting in client programs so that echo is properly disabled on Windows when
stdin is not the terminal (Matthew Stickney)
Appendix E. Release Notes

- Further fix mis-quoting of values for list-valued GUC variables in dumps (Tom Lane)
  The previous fix for quoting of `search_path` and other list-valued variables in `pg_dump` output turned out to misbehave for empty-string list elements, and it risked truncation of long file paths.

- Fix `pg_dump`’s failure to dump `REPLICA IDENTITY` properties for constraint indexes (Tom Lane)
  Manually created unique indexes were properly marked, but not those created by declaring `UNIQUE` or `PRIMARY KEY` constraints.

- Make `pg_upgrade` check that the old server was shut down cleanly (Bruce Momjian)
  The previous check could be fooled by an immediate-mode shutdown.

- Fix `contrib/hstore_plperl` to look through Perl scalar references, and to not crash if it doesn’t find a hash reference where it expects one (Tom Lane)

- Fix crash in `contrib/ltree`’s `lca()` function when the input array is empty (Pierre Ducroquet)

- Fix various error-handling code paths in which an incorrect error code might be reported (Michael Paquier, Tom Lane, Magnus Hagander)

- Rearrange makefiles to ensure that programs link to freshly-built libraries (such as `libpq.so`) rather than ones that might exist in the system library directories (Tom Lane)
  This avoids problems when building on platforms that supply old copies of PostgreSQL libraries.

- Update time zone data files to `tzdata` release 2018e for DST law changes in North Korea, plus historical corrections for Czechoslovakia.
  This update includes a redefinition of “daylight savings” in Ireland, as well as for some past years in Namibia and Czechoslovakia. In those jurisdictions, legally standard time is observed in summer, and daylight savings time in winter, so that the daylight savings offset is one hour behind standard time not one hour ahead. This does not affect either the actual UTC offset or the timezone abbreviations in use; the only known effect is that the `is_dst` column in the `pg_timezone_names` view will now be true in winter and false in summer in these cases.

E.9. Release 9.6.9

Release date: 2018-05-10

This release contains a variety of fixes from 9.6.8. For information about new features in the 9.6 major release, see Section E.18.

E.9.1. Migration to Version 9.6.9

A dump/restore is not required for those running 9.6.X.

However, if you use the `adminpack` extension, you should update it as per the first changelog entry below.

Also, if the function marking mistakes mentioned in the second and third changelog entries below affect you, you will want to take steps to correct your database catalogs.

Also, if you are upgrading from a version earlier than 9.6.8, see Section E.10.
E.9.2. Changes

• Remove public execute privilege from contrib/adminpack's pg_logfile_rotate() function (Stephen Frost)

pg_logfile_rotate() is a deprecated wrapper for the core function pg_rotate_logfile(). When that function was changed to rely on SQL privileges for access control rather than a hard-coded superuser check, pg_logfile_rotate() should have been updated as well, but the need for this was missed. Hence, if adminpack is installed, any user could request a logfile rotation, creating a minor security issue.

After installing this update, administrators should update adminpack by performing ALTER EXTENSION adminpack UPDATE in each database in which adminpack is installed. (CVE-2018-1115)

• Fix incorrect volatility markings on a few built-in functions (Thomas Munro, Tom Lane)

The functions query_to_xml, cursor_to_xml, cursor_to_xmlschema, query_to_xmlschema, and query_to_xml_and_xmlschema should be marked volatile because they execute user-supplied queries that might contain volatile operations. They were not, leading to a risk of incorrect query optimization. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard, but in case of trouble, it can be fixed by manually updating these functions' pg_proc entries, for example ALTER FUNCTION pg_catalog.query_to_xml (text, boolean, boolean, text) VOLATILE. (Note that will need to be done in each database of the installation.) Another option is to pg_upgrade the database to a version containing the corrected initial data.

• Fix incorrect parallel-safety markings on a few built-in functions (Thomas Munro, Tom Lane)

The functions brin_summarize_new_values, gin_clean_pending_list, cursor_to_xml, cursor_to_xmlschema, ts_rewrite, ts_stat, and binary_upgrade_create_empty_extension should be marked parallel-unsafe; some because they perform database modifications directly, and others because they execute user-supplied queries that might do so. They were marked parallel-restricted instead, leading to a risk of unexpected query errors. This has been repaired for new installations by correcting the initial catalog data, but existing installations will continue to contain the incorrect markings. Practical use of these functions seems to pose little hazard unless force_parallel_mode is turned on. In case of trouble, it can be fixed by manually updating these functions’ pg_proc entries, for example ALTER FUNCTION pg_catalog.brin_summarize_new_values(regclass) PARALLEL UNSAFE. (Note that will need to be done in each database of the installation.) Another option is to pg_upgrade the database to a version containing the corrected initial data.

• Avoid re-using TOAST value OIDs that match dead-but-not-yet-vacuumed TOAST entries (Pavan Deolasee)

Once the OID counter has wrapped around, it’s possible to assign a TOAST value whose OID matches a previously deleted entry in the same TOAST table. If that entry were not yet vacuumed away, this resulted in “unexpected chunk number 0 (expected 1) for toast value nnnnn” errors, which would persist until the dead entry was removed by VACUUM. Fix by not selecting such OIDs when creating a new TOAST entry.

• Change ANALYZE's algorithm for updating pg_class.reltuples (David Gould)

Previously, pages not actually scanned by ANALYZE were assumed to retain their old tuple density. In a large table where ANALYZE samples only a small fraction of the pages, this meant that the
overall tuple density estimate could not change very much, so that `reltuples` would change nearly proportionally to changes in the table’s physical size (`relpages`) regardless of what was actually happening in the table. This has been observed to result in `reltuples` becoming so much larger than reality as to effectively shut off autovacuuming. To fix, assume that `ANALYZE`’s sample is a statistically unbiased sample of the table (as it should be), and just extrapolate the density observed within those pages to the whole table.

- Avoid deadlocks in concurrent `CREATE INDEX CONCURRENTLY` commands that are run under `SERIALIZABLE` or `REPEATABLE READ` transaction isolation (Tom Lane)
- Fix possible slow execution of `REFRESH MATERIALIZED VIEW CONCURRENTLY` (Thomas Munro)
- Fix `UPDATE/DELETE ... WHERE CURRENT OF` to not fail when the referenced cursor uses an index-only-scan plan (Yugo Nagata, Tom Lane)
- Fix incorrect planning of join clauses pushed into parameterized paths (Andrew Gierth, Tom Lane)
  This error could result in misclassifying a condition as a “join filter” for an outer join when it should be a plain “filter” condition, leading to incorrect join output.
- Fix possibly incorrect generation of an index-only-scan plan when the same table column appears in multiple index columns, and only some of those index columns use operator classes that can return the column value (Kyotaro Horiguchi)
- Fix misoptimization of `CHECK` constraints having provably-NULL subclauses of top-level `AND/OR` conditions (Tom Lane, Dean Rasheed)
  This could, for example, allow constraint exclusion to exclude a child table that should not be excluded from a query.
- Fix executor crash due to double free in some `GROUPING SET` usages (Peter Geoghegan)
- Avoid crash if a table rewrite event trigger is added concurrently with a command that could call such a trigger (Álvaro Herrera, Andrew Gierth, Tom Lane)
- Avoid failure if a query-cancel or session-termination interrupt occurs while committing a prepared transaction (Stas Kelvich)
- Fix query-lifespan memory leakage in repeatedly executed hash joins (Tom Lane)
- Fix possible leak or double free of visibility map buffer pins (Amit Kapila)
- Avoid spuriously marking pages as all-visible (Dan Wood, Pavan Deolasee, Álvaro Herrera)
  This could happen if some tuples were locked (but not deleted). While queries would still function correctly, vacuum would normally ignore such pages, with the long-term effect that the tuples were never frozen. In recent releases this would eventually result in errors such as “found multixact `nnnnn` from before relminmxid `nnnnn`”.
- Fix overly strict sanity check in `heap_prepare_freeze_tuple` (Álvaro Herrera)
  This could result in incorrect “cannot freeze committed xmax” failures in databases that have been `pg_upgrade`d from 9.2 or earlier.
- Prevent dangling-pointer dereference when a C-coded before-update row trigger returns the “old” tuple (Rushabh Lathia)
- Reduce locking during autovacuum worker scheduling (Jeff Janes)
  The previous behavior caused drastic loss of potential worker concurrency in databases with many tables.
• Ensure client hostname is copied while copying `pg_stat_activity` data to local memory (Edmund Horner)
  Previously the supposedly-local snapshot contained a pointer into shared memory, allowing the clien
t hostname column to change unexpectedly if any existing session disconnected.
• Fix incorrect processing of multiple compound affixes in `ispell` dictionaries (Arthur Zakirov)
• Fix collation-aware searches (that is, indexscans using inequality operators) in SP-GiST indexes on
text columns (Tom Lane)
  Such searches would return the wrong set of rows in most non-C locales.
• Prevent query-lifespan memory leakage with SP-GiST operator classes that use traversal values
  (Anton Dignös)
• Count the number of index tuples correctly during initial build of an SP-GiST index (Tomas Vondra)
  Previously, the tuple count was reported to be the same as that of the underlying table, which is
wrong if the index is partial.
• Count the number of index tuples correctly during vacuuming of a GiST index (Andrey Borodin)
  Previously it reported the estimated number of heap tuples, which might be inaccurate, and is
certainly wrong if the index is partial.
• Fix a corner case where a streaming standby gets stuck at a WAL continuation record (Kyotaro
  Horiguchi)
• In logical decoding, avoid possible double processing of WAL data when a walsender restarts
  (Craig Ringer)
• Allow `scalarltsel` and `scalargtsel` to be used on non-core datatypes (Tomas Vondra)
• Reduce libpq’s memory consumption when a server error is reported after a large amount of query
  output has been collected (Tom Lane)
  Discard the previous output before, not after, processing the error message. On some platforms,
notably Linux, this can make a difference in the application’s subsequent memory footprint.
• Fix double-free crashes in `ecpg` (Patrick Krecker, Jeevan Ladhe)
• Fix `ecpg` to handle `long long int` variables correctly in MSVC builds (Michael Meskes, Andrew
  Gierth)
• Fix mis-quoting of values for list-valued GUC variables in dumps (Michael Paquier, Tom Lane)
  The `local_preload_libraries`, `session_preload_libraries`, `shared_preload_libraries`, and `temp_tablespaces`
variables were not correctly quoted in `pg_dump` output. This would cause problems if settings for these
variables appeared in `CREATE FUNCTION ... SET` or `ALTER DATABASE/ROLE ... SET` clauses.
• Fix `pg_recvlogical` to not fail against pre-v10 PostgreSQL servers (Michael Paquier)
  A previous fix caused `pg_recvlogical` to issue a command regardless of server version, but it should
only be issued to v10 and later servers.
• Ensure that `pg_rewind` deletes files on the target server if they are deleted from the source server
during the run (Takayuki Tsunakawa)
  Failure to do this could result in data inconsistency on the target, particularly if the file in question
is a WAL segment.
• Fix `pg_rewind` to handle tables in non-default tablespaces correctly (Takayuki Tsunakawa)
• Fix overflow handling in PL/pgSQL integer `FOR` loops (Tom Lane)
The previous coding failed to detect overflow of the loop variable on some non-gcc compilers, leading to an infinite loop.

- Adjust PL/Python regression tests to pass under Python 3.7 (Peter Eisentraut)
- Support testing PL/Python and related modules when building with Python 3 and MSVC (Andrew Dunstan)
- Fix errors in initial build of contrib/bloom indexes (Tomas Vondra, Tom Lane)
  Fix possible omission of the table’s last tuple from the index. Count the number of index tuples correctly, in case it is a partial index.
- Rename internal `b64_encode` and `b64_decode` functions to avoid conflict with Solaris 11.4 built-in functions (Rainer Orth)
- Sync our copy of the timezone library with IANA tzcode release 2018e (Tom Lane)
  This fixes the zic timezone data compiler to cope with negative daylight-savings offsets. While the PostgreSQL project will not immediately ship such timezone data, zic might be used with timezone data obtained directly from IANA, so it seems prudent to update zic now.
- Update time zone data files to tzdata release 2018d for DST law changes in Palestine and Antarctica (Casey Station), plus historical corrections for Portugal and its colonies, as well as Enderbury, Jamaica, Turks & Caicos Islands, and Uruguay.

**E.10. Release 9.6.8**

**Release date:** 2018-03-01

This release contains a variety of fixes from 9.6.7. For information about new features in the 9.6 major release, see Section E.18.

**E.10.1. Migration to Version 9.6.8**

A dump/restore is not required for those running 9.6.X.

However, if you run an installation in which not all users are mutually trusting, or if you maintain an application or extension that is intended for use in arbitrary situations, it is strongly recommended that you read the documentation changes described in the first changelog entry below, and take suitable steps to ensure that your installation or code is secure.

Also, the changes described in the second changelog entry below may cause functions used in index expressions or materialized views to fail during auto-analyze, or when reloading from a dump. After upgrading, monitor the server logs for such problems, and fix affected functions.

Also, if you are upgrading from a version earlier than 9.6.7, see Section E.11.

**E.10.2. Changes**

- Document how to configure installations and applications to guard against search-path-dependent trojan-horse attacks from other users (Noah Misch)

  Using a `search_path` setting that includes any schemas writable by a hostile user enables that user to capture control of queries and then run arbitrary SQL code with the permissions of the
Appendix E. Release Notes

attacked user. While it is possible to write queries that are proof against such hijacking, it is notationally tedious, and it’s very easy to overlook holes. Therefore, we now recommend configurations in which no untrusted schemas appear in one’s search path. Relevant documentation appears in Section 5.8.6 (for database administrators and users), Section 32.1 (for application authors), Section 36.15.1 (for extension authors), and CREATE FUNCTION (for authors of SECURITY DEFINER functions). (CVE-2018-1058)

- Avoid use of insecure `search_path` settings in `pg_dump` and other client programs (Noah Misch, Tom Lane)

`pg_dump`, `pg_upgrade`, `vacuumdb` and other PostgreSQL-provided applications were themselves vulnerable to the type of hijacking described in the previous changelog entry; since these applications are commonly run by superusers, they present particularly attractive targets. To make them secure whether or not the installation as a whole has been secured, modify them to include only the `pg_catalog` schema in their `search_path` settings. Autovacuum worker processes now do the same, as well.

In cases where user-provided functions are indirectly executed by these programs — for example, user-provided functions in index expressions — the tighter `search_path` may result in errors, which will need to be corrected by adjusting those user-provided functions to not assume anything about what search path they are invoked under. That has always been good practice, but now it will be necessary for correct behavior. (CVE-2018-1058)

- Fix misbehavior of concurrent-update rechecks with CTE references appearing in subplans (Tom Lane)

If a CTE (`WITH` clause reference) is used in an InitPlan or SubPlan, and the query requires a recheck due to trying to update or lock a concurrently-updated row, incorrect results could be obtained.

- Fix planner failures with overlapping mergejoin clauses in an outer join (Tom Lane)

These mistakes led to “left and right pathkeys do not match in mergejoin” or “outer pathkeys do not match mergeclauses” planner errors in corner cases.

- Repair `pg_upgrade`’s failure to preserve `relfrozenxid` for materialized views (Tom Lane, Andres Freund)

This oversight could lead to data corruption in materialized views after an upgrade, manifesting as “could not access status of transaction” or “found xmin from before relfrozenxid” errors. The problem would be more likely to occur in seldom-refreshed materialized views, or ones that were maintained only with `REFRESH MATERIALIZED VIEW CONCURRENTLY`.

If such corruption is observed, it can be repaired by refreshing the materialized view (without `CONCURRENTLY`).

- Fix incorrect reporting of PL/Python function names in error `CONTEXT` stacks (Tom Lane)

An error occurring within a nested PL/Python function call (that is, one reached via a SPI query from another PL/Python function) would result in a stack trace showing the inner function’s name twice, rather than the expected results. Also, an error in a nested PL/Python `DO` block could result in a null pointer dereference crash on some platforms.

- Allow contrib/auto_explain’s `log_min_duration` setting to range up to `INT_MAX`, or about 24 days instead of 35 minutes (Tom Lane)

- Mark assorted GUC variables as `PGDLLIMPORT`, to ease porting extension modules to Windows (Metin Doslu)

Release date: 2018-02-08

This release contains a variety of fixes from 9.6.6. For information about new features in the 9.6 major release, see Section E.18.

E.11.1. Migration to Version 9.6.7

A dump/restore is not required for those running 9.6.X.

However, if you use contrib/cube’s `~>` operator, see the entry below about that.

Also, if you are upgrading from a version earlier than 9.6.6, see Section E.12.

E.11.2. Changes

• Ensure that all temporary files made by pg_upgrade are non-world-readable (Tom Lane, Noah Misch)

  pg_upgrade normally restricts its temporary files to be readable and writable only by the calling user. But the temporary file containing `pg_dumpall -g` output would be group- or world-readable, or even writable, if the user’s `umask` setting allows. In typical usage on multi-user machines, the `umask` and/or the working directory’s permissions would be tight enough to prevent problems; but there may be people using pg_upgrade in scenarios where this oversight would permit disclosure of database passwords to unfriendly eyes. (CVE-2018-1053)

• Fix vacuuming of tuples that were updated while key-share locked (Andres Freund, Álvaro Herrera)

  In some cases `VACUUM` would fail to remove such tuples even though they are now dead, leading to assorted data corruption scenarios.

• Ensure that vacuum will always clean up the pending-insertions list of a GIN index (Masahiko Sawada)

  This is necessary to ensure that dead index entries get removed. The old code got it backwards, allowing vacuum to skip the cleanup if some other process were running cleanup concurrently, thus risking invalid entries being left behind in the index.

• Fix inadequate buffer locking in some LSN fetches (Jacob Champion, Asim Praveen, Ashwin Agrawal)

  These errors could result in misbehavior under concurrent load. The potential consequences have not been characterized fully.

• Fix incorrect query results from cases involving flattening of subqueries whose outputs are used in `GROUPING SETS` (Heikki Linnakangas)

• Avoid unnecessary failure in a query on an inheritance tree that occurs concurrently with some child table being removed from the tree by `ALTER TABLE NO INHERIT` (Tom Lane)

• Fix spurious deadlock failures when multiple sessions are running `CREATE INDEX CONCURRENTLY` (Jeff Janes)

• Fix failures when an inheritance tree contains foreign child tables (Etsuro Fujita)
A mix of regular and foreign tables in an inheritance tree resulted in creation of incorrect plans for `UPDATE` and `DELETE` queries. This led to visible failures in some cases, notably when there are row-level triggers on a foreign child table.

- Repair failure with correlated sub-SELECT inside VALUES inside a LATERAL subquery (Tom Lane)
- Fix “could not devise a query plan for the given query” planner failure for some cases involving nested `UNION ALL` inside a lateral subquery (Tom Lane)
- Fix logical decoding to correctly clean up disk files for crashed transactions (Atsushi Torikoshi)
  Logical decoding may spill WAL records to disk for transactions generating many WAL records. Normally these files are cleaned up after the transaction’s commit or abort record arrives; but if no such record is ever seen, the removal code misbehaved.
- Fix walsender timeout failure and failure to respond to interrupts when processing a large transaction (Petr Jelinek)
- Fix `has_sequence_privilege()` to support WITH GRANT OPTION tests, as other privilege-testing functions do (Joe Conway)
- In databases using UTF8 encoding, ignore any XML declaration that asserts a different encoding (Pavel Stehule, Noah Misch)
  We always store XML strings in the database encoding, so allowing libxml to act on a declaration of another encoding gave wrong results. In encodings other than UTF8, we don’t promise to support non-ASCII XML data anyway, so retain the previous behavior for bug compatibility. This change affects only `xpath()` and related functions; other XML code paths already acted this way.
- Provide for forward compatibility with future minor protocol versions (Robert Haas, Badrul Chowdhury)
  Up to now, PostgreSQL servers simply rejected requests to use protocol versions newer than 3.0, so that there was no functional difference between the major and minor parts of the protocol version number. Allow clients to request versions 3.x without failing, sending back a message showing that the server only understands 3.0. This makes no difference at the moment, but back-patching this change should allow speedier introduction of future minor protocol upgrades.
- Cope with failure to start a parallel worker process (Amit Kapila, Robert Haas)
  Parallel query previously tended to hang indefinitely if a worker could not be started, as the result of `fork()` failure or other low-probability problems.
- Fix collection of EXPLAIN statistics from parallel workers (Amit Kapila, Thomas Munro)
- Avoid unsafe alignment assumptions when working with `__int128` (Tom Lane)
  Typically, compilers assume that `__int128` variables are aligned on 16-byte boundaries, but our memory allocation infrastructure isn’t prepared to guarantee that, and increasing the setting of `MAXALIGN` seems infeasible for multiple reasons. Adjust the code to allow use of `__int128` only when we can tell the compiler to assume lesser alignment. The only known symptom of this problem so far is crashes in some parallel aggregation queries.
- Prevent stack-overflow crashes when planning extremely deeply nested set operations (UNION/INTERSECT/EXCEPT) (Tom Lane)
- Fix null-pointer crashes for some types of LDAP URLs appearing in `pg_hba.conf` (Thomas Munro)
- Fix sample `INSTR()` functions in the PL/pgSQL documentation (Yugo Nagata, Tom Lane)
  These functions are stated to be Oracle® compatible, but they weren’t exactly. In particular, there was a discrepancy in the interpretation of a negative third parameter: Oracle thinks that a negative
Appendix E. Release Notes

value indicates the last place where the target substring can begin, whereas our functions took it as the last place where the target can end. Also, Oracle throws an error for a zero or negative fourth parameter, whereas our functions returned zero.

The sample code has been adjusted to match Oracle’s behavior more precisely. Users who have copied this code into their applications may wish to update their copies.

• Fix pg_dump to make ACL (permissions), comment, and security label entries reliably identifiable in archive output formats (Tom Lane)

The “tag” portion of an ACL archive entry was usually just the name of the associated object. Make it start with the object type instead, bringing ACLs into line with the convention already used for comment and security label archive entries. Also, fix the comment and security label entries for the whole database, if present, to make their tags start with DATABASE so that they also follow this convention. This prevents false matches in code that tries to identify large-object-related entries by seeing if the tag starts with LARGE OBJECT. That could have resulted in misclassifying entries as data rather than schema, with undesirable results in a schema-only or data-only dump.

Note that this change has user-visible results in the output of pg_restore --list.

• Rename pg_rewind’s copy_file_range function to avoid conflict with new Linux system call of that name (Andres Freund)

This change prevents build failures with newer glibc versions.

• In ecpg, detect indicator arrays that do not have the correct length and report an error (David Rader)

• Change the behavior of contrib/cube’s cube ~> int operator to make it compatible with KNN search (Alexander Korotkov)

The meaning of the second argument (the dimension selector) has been changed to make it predictable which value is selected even when dealing with cubes of varying dimensionalities.

This is an incompatible change, but since the point of the operator was to be used in KNN searches, it seems rather useless as-is. After installing this update, any expression indexes or materialized views using this operator will need to be reindexed/refreshed.

• Avoid triggering a libc assertion in contrib/hstore, due to use of memcpy() with equal source and destination pointers (Tomas Vondra)

• Fix incorrect display of tuples’ null bitmaps in contrib/pageinspect (Maksim Milyutin)

• In contrib/postgres_fdw, avoid “outer pathkeys do not match mergeclauses” planner error when constructing a plan involving a remote join (Robert Haas)

• Provide modern examples of how to auto-start Postgres on macOS (Tom Lane)

The scripts in contrib/start-scripts/osx use infrastructure that’s been deprecated for over a decade, and which no longer works at all in macOS releases of the last couple of years. Add a new subdirectory contrib/start-scripts/macos containing scripts that use the newer launchd infrastructure.

• Fix incorrect selection of configuration-specific libraries for OpenSSL on Windows (Andrew Dunstan)

• Support linking to MinGW-built versions of libperl (Noah Misch)

This allows building PL/Perl with some common Perl distributions for Windows.

• Fix MSVC build to test whether 32-bit libperl needs -D_USE_32BIT_TIME_T (Noah Misch)

Available Perl distributions are inconsistent about what they expect, and lack any reliable means of reporting it, so resort to a build-time test on what the library being used actually does.
• On Windows, install the crash dump handler earlier in postmaster startup (Takayuki Tsunakawa)
  This may allow collection of a core dump for some early-startup failures that did not produce a
dump before.
• On Windows, avoid encoding-conversion-related crashes when emitting messages very early in
  postmaster startup (Takayuki Tsunakawa)
• Use our existing Motorola 68K spinlock code on OpenBSD as well as NetBSD (David Carlier)
• Add support for spinlocks on Motorola 88K (David Carlier)
• Update time zone data files to tzdata release 2018c for DST law changes in Brazil, Sao Tome and
  Principe, plus historical corrections for Bolivia, Japan, and South Sudan. The US/Pacific-New
  zone has been removed (it was only an alias for America/Los_Angeles anyway).

E.12. Release 9.6.6

Release date: 2017-11-09
This release contains a variety of fixes from 9.6.5. For information about new features in the 9.6 major
release, see Section E.18.

E.12.1. Migration to Version 9.6.6
A dump/restore is not required for those running 9.6.X.
However, if you use BRIN indexes, see the fourth changelog entry below.
Also, if you are upgrading from a version earlier than 9.6.4, see Section E.14.

E.12.2. Changes

• Ensure that INSERT ... ON CONFLICT DO UPDATE checks table permissions and RLS policies
  in all cases (Dean Rasheed)
  The update path of INSERT ... ON CONFLICT DO UPDATE requires SELECT permission on the
  columns of the arbiter index, but it failed to check for that in the case of an arbiter specified by
  constraint name. In addition, for a table with row level security enabled, it failed to check updated
  rows against the table’s SELECT policies (regardless of how the arbiter index was specified). (CVE-
  2017-15099)
• Fix crash due to rowtype mismatch in jsonb_populate_recordset() (Michael Paquier, Tom
  Lane)
  These functions used the result rowtype specified in the FROM ... AS clause without checking
  that it matched the actual rowtype of the supplied tuple value. If it didn’t, that would usually result
  in a crash, though disclosure of server memory contents seems possible as well. (CVE-2017-15098)
• Fix sample server-start scripts to become $PGUSER before opening $PGLOG (Noah Misch)
  Previously, the postmaster log file was opened while still running as root. The database owner
could therefore mount an attack against another system user by making $PGLOG be a symbolic link
to some other file, which would then become corrupted by appending log messages.
By default, these scripts are not installed anywhere. Users who have made use of them will need to manually recopy them, or apply the same changes to their modified versions. If the existing $PGLOG file is root-owned, it will need to be removed or renamed out of the way before restarting the server with the corrected script. (CVE-2017-12172)

- Fix BRIN index summarization to handle concurrent table extension correctly (Álvaro Herrera)
  Previously, a race condition allowed some table rows to be omitted from the index. It may be necessary to reindex existing BRIN indexes to recover from past occurrences of this problem.

- Fix possible failures during concurrent updates of a BRIN index (Tom Lane)
  These race conditions could result in errors like “invalid index offnum” or “inconsistent range map”.

- Fix crash when logical decoding is invoked from a SPI-using function, in particular any function written in a PL language (Tom Lane)

- Fix incorrect query results when multiple GROUPING SETS columns contain the same simple variable (Tom Lane)

- Fix incorrect parallelization decisions for nested queries (Amit Kapila, Kuntal Ghosh)

- Fix parallel query handling to not fail when a recently-used role is dropped (Amit Kapila)

- Fix json_build_array(), json_build_object(), and their jsonb equivalents to handle explicit VARIADIC arguments correctly (Michael Paquier)

- Properly reject attempts to convert infinite float values to type numeric (Tom Lane, KaiGai Kohei)
  Previously the behavior was platform-dependent.

- Fix corner-case crashes when columns have been added to the end of a view (Tom Lane)

- Record proper dependencies when a view or rule contains FieldSelect or FieldStore expression nodes (Tom Lane)
  Lack of these dependencies could allow a column or data type DROP to go through when it ought to fail, thereby causing later uses of the view or rule to get errors. This patch does not do anything to protect existing views/rules, only ones created in the future.

- Correctly detect hashability of range data types (Tom Lane)
  The planner mistakenly assumed that any range type could be hashed for use in hash joins or hash aggregation, but actually it must check whether the range’s subtype has hash support. This does not affect any of the built-in range types, since they’re all hashable anyway.

- Correctly ignore RelabelType expression nodes when determining relation distinctness (David Rowley)
  This allows the intended optimization to occur when a subquery has a result column of type varchar.

- Prevent sharing transition states between ordered-set aggregates (David Rowley)
  This causes a crash with the built-in ordered-set aggregates, and probably with user-written ones as well. v11 and later will include provisions for dealing with such cases safely, but in released branches, just disable the optimization.

- Prevent idle_in_transaction_session_timeout from being ignored when a statement_timeout occurred earlier (Lukas Fittl)

- Fix low-probability loss of NOTIFY messages due to XID wraparound (Marko Tiikkaja, Tom Lane)
  If a session executed no queries, but merely listened for notifications, for more than 2 billion transactions, it started to miss some notifications from concurrently-committing transactions.
• Avoid SIGBUS crash on Linux when a DSM memory request exceeds the space available in tmpfs (Thomas Munro)

• Reduce the frequency of data flush requests during bulk file copies to avoid performance problems on macOS, particularly with its new APFS file system (Tom Lane)

• Prevent low-probability crash in processing of nested trigger firings (Tom Lane)

• Allow copy's FREEZE option to work when the transaction isolation level is REPEATABLE READ or higher (Noah Misch)

  This case was unintentionally broken by a previous bug fix.

• Correctly restore the umask setting when file creation fails in copy or lo_export() (Peter Eisen-traut)

• Give a better error message for duplicate column names in ANALYZE (Nathan Bossart)

• Add missing cases in GetCommandLogLevel(), preventing errors when certain SQL commands are used while log_statement is set to ddl (Michael Paquier)

• Fix mis-parsing of the last line in a non-newline-terminated pg_hba.conf file (Tom Lane)

• Fix AggGetAggref() to return the correct Aggref nodes to aggregate final functions whose transition calculations have been merged (Tom Lane)

• Fix pg_dump to ensure that it emits GRANT commands in a valid order (Stephen Frost)

• Fix pg_basebackup’s matching of tablespace paths to canonicalize both paths before comparing (Michael Paquier)

  This is particularly helpful on Windows.

• Fix libpq to not require user’s home directory to exist (Tom Lane)

  In v10, failure to find the home directory while trying to read ~/.pgpass was treated as a hard error, but it should just cause that file to not be found. Both v10 and previous release branches made the same mistake when reading ~/.pg_service.conf, though this was less obvious since that file is not sought unless a service name is specified.

• Fix libpq to guard against integer overflow in the row count of a PGresult (Michael Paquier)

• Fix ecpg’s handling of out-of-scope cursor declarations with pointer or array variables (Michael Meskes)

  In ecpglib, correctly handle backslashes in string literals depending on whether standard_conforming_strings is set (Tsunakawa Takayuki)

  Make ecpglib’s Informix-compatibility mode ignore fractional digits in integer input strings, as expected (Gao Zengqi, Michael Meskes)

• Fix ecpg’s regression tests to work reliably on Windows (Christian Ullrich, Michael Meskes)

• Fix missing temp-install prerequisites for check-like Make targets (Noah Misch)

  Some non-default test procedures that are meant to work like make check failed to ensure that the temporary installation was up to date.

• Sync our copy of the timezone library with IANA release tzcode2017c (Tom Lane)

  This fixes various issues; the only one likely to be user-visible is that the default DST rules for a POSIX-style zone name, if no posixrules file exists in the timezone data directory, now match current US law rather than what it was a dozen years ago.
• Update time zone data files to tzdata release 2017c for DST law changes in Fiji, Namibia, Northern Cyprus, Sudan, Tonga, and Turks & Caicos Islands, plus historical corrections for Alaska, Apia, Burma, Calcutta, Detroit, Ireland, Namibia, and Pago Pago.

E.13. Release 9.6.5

Release date: 2017-08-31

This release contains a small number of fixes from 9.6.4. For information about new features in the 9.6 major release, see Section E.18.

E.13.1. Migration to Version 9.6.5

A dump/restore is not required for those running 9.6.X.

However, if you are upgrading from a version earlier than 9.6.4, see Section E.14.

E.13.2. Changes

• Show foreign tables in information_schema.table_privileges view (Peter Eisentraut)

  All other relevant information_schema views include foreign tables, but this one ignored them.

  Since this view definition is installed by initdb, merely upgrading will not fix the problem. If you need to fix this in an existing installation, you can, as a superuser, do this in psql:

        SET search_path TO information_schema;
        CREATE OR REPLACE VIEW table_privileges AS
            SELECT CAST(u_grantor.rolname AS sql_identifier) AS grantor,
                   CAST(grantee.rolname AS sql_identifier) AS grantee,
                   CAST(current_database() AS sql_identifier) AS table_catalog,
                   CAST(nc.nspname AS sql_identifier) AS table_schema,
                   CAST(c.relname AS sql_identifier) AS table_name,
                   CAST(c.prtype AS character_data) AS privilege_type,
                   CAST(
                       CASE WHEN
                           -- object owner always has grant options
                           pg_has_role(grantee.oid, c.relowner, 'USAGE')
                           OR c.grantable
                       THEN 'YES' ELSE 'NO' END AS yes_or_no) AS is_grantable,
                   CAST(CASE WHEN c.prtype = 'SELECT' THEN 'YES' ELSE 'NO' END AS yes_or_no) AS with_hierarchy
            FROM (SELECT oid, relname, relnamespace, relkind, relowner, (aclexplode(coalesce(relacl, acldefault('r', relowner)))).* FROM pg_class AS c (oid, relname, relnamespace, relkind, relowner, prtype, nspname, pg_namespace nc, pg_authid u_grantor,
            (SELECT oid, rolname FROM pg_authid
            UNION ALL
            SELECT 0::oid, 'PUBLIC')
            ) AS grantee (oid, rolname)
WHERE c.relnamespace = nc.oid
AND c.relkind IN ('r', 'v', 'f')
AND c.grantee = grantee.oid
AND c.grantor = u_grantor.oid
AND c.prtype IN ('INSERT', 'SELECT', 'UPDATE', 'DELETE', 'TRUNCATE', 'REFERENCES', 'TRIGGER')
AND (pg_has_role(u_grantor.oid, 'USAGE')
OR pg_has_role(grantee.oid, 'USAGE')
OR grantee.rolname = 'PUBLIC');

This must be repeated in each database to be fixed, including template0.

- Clean up handling of a fatal exit (e.g., due to receipt of SIGTERM) that occurs while trying to execute a ROLLBACK of a failed transaction (Tom Lane)

  This situation could result in an assertion failure. In production builds, the exit would still occur, but it would log an unexpected message about “cannot drop active portal”.

- Remove assertion that could trigger during a fatal exit (Tom Lane)

- Correctly identify columns that are of a range type or domain type over a composite type or domain type being searched for (Tom Lane)

  Certain ALTER commands that change the definition of a composite type or domain type are supposed to fail if there are any stored values of that type in the database, because they lack the infrastructure needed to update or check such values. Previously, these checks could miss relevant values that are wrapped inside range types or sub-domains, possibly allowing the database to become inconsistent.

- Prevent crash when passing fixed-length pass-by-reference data types to parallel worker processes (Tom Lane)

- Fix crash in pg_restore when using parallel mode and using a list file to select a subset of items to restore (Fabrizio de Royes Mello)

- Change ecpg’s parser to allow RETURNING clauses without attached C variables (Michael Meskes)

  This allows ecpg programs to contain SQL constructs that use RETURNING internally (for example, inside a CTE) rather than using it to define values to be returned to the client.

- Change ecpg’s parser to recognize backslash continuation of C preprocessor command lines (Michael Meskes)

- Improve selection of compiler flags for PL/Perl on Windows (Tom Lane)

  This fix avoids possible crashes of PL/Perl due to inconsistent assumptions about the width of time_t values. A side-effect that may be visible to extension developers is that _USE_32BIT_TIME_T is no longer defined globally in PostgreSQL Windows builds. This is not expected to cause problems, because type time_t is not used in any PostgreSQL API definitions.

- Fix make check to behave correctly when invoked via a non-GNU make program (Thomas Munro)

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**E.14. Release 9.6.4**

**Release date:** 2017-08-10

This release contains a variety of fixes from 9.6.3. For information about new features in the 9.6 major release, see Section E.18.
E.14.1. Migration to Version 9.6.4

A dump/restore is not required for those running 9.6.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are upgrading from a version earlier than 9.6.3, see Section E.15.

E.14.2. Changes

• Further restrict visibility of `pg_user_mappings.umoptions`, to protect passwords stored as user mapping options (Noah Misch)

The fix for CVE-2017-7486 was incorrect: it allowed a user to see the options in her own user mapping, even if she did not have `USAGE` permission on the associated foreign server. Such options might include a password that had been provided by the server owner rather than the user herself. Since `information_schema.user_mapping_options` does not show the options in such cases, `pg_user_mappings` should not either. (CVE-2017-7547)

By itself, this patch will only fix the behavior in newly initdb’d databases. If you wish to apply this change in an existing database, you will need to do the following:

1. Restart the postmaster after adding `allow_system_table_mods = true` to `postgresql.conf`. (In versions supporting `ALTER SYSTEM`, you can use that to make the configuration change, but you’ll still need a restart.)

2. In each database of the cluster, run the following commands as superuser:

   ```sql
   SET search_path = pg_catalog;
   CREATE OR REPLACE VIEW pg_user_mappings AS
   SELECT
       U.oid AS umid,
       S.oid AS srvid,
       S.srvname AS srvname,
       U.umuser AS umuser,
       CASE WHEN U.umuser = 0 THEN
           'public'
       ELSE
           A.rolname
       END AS usename,
       CASE WHEN U.umuser = 0 AND pg_has_role(S.srvowner, 'USAGE')
           OR (SELECT rolsuper FROM pg_authid WHERE rolname = current_user)
           THEN U.umoptions
       ELSE NULL END AS umoptions
   FROM pg_user_mapping U
   LEFT JOIN pg_authid A ON (A.oid = U.umuser) JOIN pg_foreign_server S ON (U.umserver = S.oid);
   ```

3. Do not forget to include the `template0` and `template1` databases, or the vulnerability will still exist in databases you create later. To fix `template0`, you’ll need to temporarily make it accept connections. In PostgreSQL 9.5 and later, you can use

   ```sql
   ALTER DATABASE template0 WITH ALLOW_CONNECTIONS true;
   ```
and then after fixing `template0`, undo that with

```
ALTER DATABASE template0 WITH ALLOW_CONNECTIONS false;
```

In prior versions, instead use

```
UPDATE pg_database SET datallowconn = true WHERE datname = 'template0';
UPDATE pg_database SET datallowconn = false WHERE datname = 'template0';
```

4. Finally, remove the `allow_system_table_mods` configuration setting, and again restart the postmaster.

- Disallow empty passwords in all password-based authentication methods (Heikki Linnakangas)

  libpq ignores empty password specifications, and does not transmit them to the server. So, if a user’s password has been set to the empty string, it’s impossible to log in with that password via `psql` or other libpq-based clients. An administrator might therefore believe that setting the password to empty is equivalent to disabling password login. However, with a modified or non-libpq-based client, logging in could be possible, depending on which authentication method is configured. In particular the most common method, `md5`, accepted empty passwords. Change the server to reject empty passwords in all cases. (CVE-2017-7546)

- Make `lo_put()` check for `UPDATE` privilege on the target large object (Tom Lane, Michael Paquier)

  `lo_put()` should surely require the same permissions as `lowrite()`, but the check was missing, allowing any user to change the data in a large object. (CVE-2017-7548)

- Correct the documentation about the process for upgrading standby servers with `pg_upgrade` (Bruce Momjian)

  The previous documentation instructed users to start/stop the primary server after running `pg_upgrade` but before syncing the standby servers. This sequence is unsafe.

- Fix concurrent locking of tuple update chains (Álvaro Herrera)

  If several sessions concurrently lock a tuple update chain with nonconflicting lock modes using an old snapshot, and they all succeed, it was possible for some of them to nonetheless fail (and conclude there is no live tuple version) due to a race condition. This had consequences such as foreign-key checks failing to see a tuple that definitely exists but is being updated concurrently.

- Fix potential data corruption when freezing a tuple whose `XMAX` is a multixact with exactly one still-interesting member (Teodor Sigaev)

- Avoid integer overflow and ensuing crash when sorting more than one billion tuples in-memory (Sergey Koposov)

- On Windows, retry process creation if we fail to reserve the address range for our shared memory in the new process (Tom Lane, Amit Kapila)

  This is expected to fix infrequent child-process-launch failures that are probably due to interference from antivirus products.

- Fix low-probability corruption of shared predicate-lock hash table in Windows builds (Thomas Munro, Tom Lane)

- Avoid logging clean closure of an SSL connection as though it were a connection reset (Michael Paquier)

- Prevent sending SSL session tickets to clients (Tom Lane)

  This fix prevents reconnection failures with ticket-aware client-side SSL code.
• Fix code for setting tcp_keepalives_idle on Solaris (Tom Lane)
• Fix statistics collector to honor inquiry messages issued just after a postmaster shutdown and immediate restart (Tom Lane)
  Statistics inquiries issued within half a second of the previous postmaster shutdown were effectively ignored.
• Ensure that the statistics collector’s receive buffer size is at least 100KB (Tom Lane)
  This reduces the risk of dropped statistics data on older platforms whose default receive buffer size is less than that.
• Fix possible creation of an invalid WAL segment when a standby is promoted just after it processes an XLOG_SWITCH WAL record (Andres Freund)
• Fix walsender to exit promptly when client requests shutdown (Tom Lane)
• Fix SIGHUP and SIGUSR1 handling in walsender processes (Petr Jelinek, Andres Freund)
• Prevent walsender-triggered panics during shutdown checkpoints (Andres Freund, Michael Paquier)
• Fix unnecessarily slow restarts of walreceiver processes due to race condition in postmaster (Tom Lane)
• Fix leakage of small subtransactions spilled to disk during logical decoding (Andres Freund)
  This resulted in temporary files consuming excessive disk space.
• Reduce the work needed to build snapshots during creation of logical-decoding slots (Andres Freund, Petr Jelinek)
  The previous algorithm was infeasibly expensive on a server with a lot of open transactions.
• Fix race condition that could indefinitely delay creation of logical-decoding slots (Andres Freund, Petr Jelinek)
• Reduce overhead in processing syscache invalidation events (Tom Lane)
  This is particularly helpful for logical decoding, which triggers frequent cache invalidation.
• Remove incorrect heuristic used in some cases to estimate join selectivity based on the presence of foreign-key constraints (David Rowley)
  In some cases where a multi-column foreign key constraint existed but did not exactly match a query’s join structure, the planner used an estimation heuristic that turns out not to work well at all. Revert such cases to the way they were estimated before 9.6.
• Fix cases where an INSERT or UPDATE assigns to more than one element of a column that is of domain-over-array type (Tom Lane)
• Allow window functions to be used in sub-SELECTs that are within the arguments of an aggregate function (Tom Lane)
• Ensure that a view’s CHECK OPTIONS clause is enforced properly when the underlying table is a foreign table (Etsuro Fujita)
  Previously, the update might get pushed entirely to the foreign server, but the need to verify the view conditions was missed if so.
• Move autogenerated array types out of the way during ALTER ... RENAME (Vik Fearing)
  Previously, we would rename a conflicting autogenerated array type out of the way during CREATE; this fix extends that behavior to renaming operations.
• Fix dangling pointer in `ALTER TABLE` when there is a comment on a constraint belonging to the
table (David Rowley)
Re-applying the comment to the reconstructed constraint could fail with a weird error message, or
even crash.

• Ensure that `ALTER USER ... SET` accepts all the syntax variants that `ALTER ROLE ... SET`
does (Peter Eisentraut)

• Allow a foreign table’s `CHECK` constraints to be initially `NOT VALID` (Amit Langote)

CREATE TABLE silently drops `NOT VALID` specifiers for `CHECK` constraints, reasoning that the
table must be empty so the constraint can be validated immediately. But this is wrong for `CREATE
FOREIGN TABLE`, where there’s no reason to suppose that the underlying table is empty, and even
if it is it’s no business of ours to decide that the constraint can be treated as valid going forward.
Skip this “optimization” for foreign tables.

• Properly update dependency info when changing a datatype I/O function’s argument or return type
from `opaque` to the correct type (Heikki Linnakangas)

CREATE TYPE updates I/O functions declared in this long-obsolete style, but it forgot to record a
dependency on the type, allowing a subsequent `DROP TYPE` to leave broken function definitions
behind.

• Allow parallelism in the query plan when `COPY` copies from a query’s result (Andres Freund)

• Reduce memory usage when `ANALYZE` processes a `tsvector` column (Heikki Linnakangas)

• Fix unnecessary precision loss and sloppy rounding when multiplying or dividing `money` values by
integers or floats (Tom Lane)

• Tighten checks for whitespace in functions that parse identifiers, such as `regprocedurein()`
(Tom Lane)

Depending on the prevailing locale, these functions could misinterpret fragments of multibyte char-
acters as whitespace.

• Use relevant `#define` symbols from Perl while compiling PL/Perl (Ashutosh Sharma, Tom Lane)
This avoids portability problems, typically manifesting as a “handshake” mismatch during library
load, when working with recent Perl versions.

• In `libpq`, reset GSS/SASL and SSPI authentication state properly after a failed connection attempt
(Michael Paquier)

Failure to do this meant that when falling back from SSL to non-SSL connections, a GSS/SASL
failure in the SSL attempt would always cause the non-SSL attempt to fail. SSPI did not fail, but it
leaked memory.

• In `psql`, fix failure when `COPY FROM STDIN` is ended with a keyboard EOF signal and then another
`COPY FROM STDIN` is attempted (Thomas Munro)

This misbehavior was observed on BSD-derived platforms (including macOS), but not on most
others.

• Fix `pg_dump` and `pg_restore` to emit `REFRESH MATERIALIZED VIEW` commands last (Tom Lane)
This prevents errors during dump/restore when a materialized view refers to tables owned by a
different user.

• Improve `pg_dump/pg_restore`’s reporting of error conditions originating in `zlib` (Vladimir Kun-
schikov, Álvaro Herrera)

• Fix `pg_dump` with the `--clean` option to drop event triggers as expected (Tom Lane)
Appendix E. Release Notes

It also now correctly assigns ownership of event triggers; before, they were restored as being owned by the superuser running the restore script.

- Fix `pg_dump` with the `--clean` option to not fail when the `public` schema doesn’t exist (Stephen Frost)
- Fix `pg_dump` to not emit invalid SQL for an empty operator class (Daniel Gustafsson)
- Fix `pg_dump` output to stdout on Windows (Kuntal Ghosh)
  A compressed plain-text dump written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.
- Fix `pg_get_ruledef()` to print correct output for the ON SELECT rule of a view whose columns have been renamed (Tom Lane)
  In some corner cases, `pg_dump` relies on `pg_get_ruledef()` to dump views, so that this error could result in dump/reload failures.
- Fix dumping of outer joins with empty constraints, such as the result of a NATURAL LEFT JOIN with no common columns (Tom Lane)
- Fix dumping of function expressions in the FROM clause in cases where the expression does not deparse into something that looks like a function call (Tom Lane)
- Fix `pg_basebackup` output to stdout on Windows (Haribabu Kommi)
  A backup written to stdout would contain corrupt data due to failure to put the file descriptor into binary mode.
- Fix `pg_rewind` to correctly handle files exceeding 2GB (Kuntal Ghosh, Michael Paquier)
  Ordinarily such files won’t appear in PostgreSQL data directories, but they could be present in some cases.
- Fix `pg_upgrade` to ensure that the ending WAL record does not have `wal_level = minimum` (Bruce Momjian)
  This condition could prevent upgraded standby servers from reconnecting.
- Fix `pg_xlogdump`’s computation of WAL record length (Andres Freund)
- In `postgres_fdw`, re-establish connections to remote servers after ALTER SERVER or ALTER USER MAPPING commands (Kyotaro Horiguchi)
  This ensures that option changes affecting connection parameters will be applied promptly.
- In `postgres_fdw`, allow cancellation of remote transaction control commands (Robert Haas, Rafia Sabih)
  This change allows us to quickly escape a wait for an unresponsive remote server in many more cases than previously.
- Increase `MAX_SYSCACHE_CALLBACKS` to provide more room for extensions (Tom Lane)
- Always use `-fPIC`, not `-fpic`, when building shared libraries with gcc (Tom Lane)
  This supports larger extension libraries on platforms where it makes a difference.
- In MSVC builds, handle the case where the openssl library is not within a `VC` subdirectory (Andrew Dunstan)
- In MSVC builds, add proper include path for libxml2 header files (Andrew Dunstan)
  This fixes a former need to move things around in standard Windows installations of libxml2.
- In MSVC builds, recognize a Tcl library that is named `tcl86.lib` (Noah Misch)
Appendix E. Release Notes

• In MSVC builds, honor PROVE_FLAGS settings on vcregress.pl's command line (Andrew Dunstan)

E.15. Release 9.6.3

Release date: 2017-05-11

This release contains a variety of fixes from 9.6.2. For information about new features in the 9.6 major release, see Section E.18.

E.15.1. Migration to Version 9.6.3

A dump/restore is not required for those running 9.6.X.

However, if you use foreign data servers that make use of user passwords for authentication, see the first changelog entry below.

Also, if you are using third-party replication tools that depend on “logical decoding”, see the fourth changelog entry below.

Also, if you are upgrading from a version earlier than 9.6.2, see Section E.16.

E.15.2. Changes

• Restrict visibility of pg_user_mappings.umoptions, to protect passwords stored as user mapping options (Michael Paquier, Feike Steenbergen)

The previous coding allowed the owner of a foreign server object, or anyone he has granted server USAGE permission to, to see the options for all user mappings associated with that server. This might well include passwords for other users. Adjust the view definition to match the behavior of information_schema.user_mapping_options, namely that these options are visible to the user being mapped, or if the mapping is for PUBLIC and the current user is the server owner, or if the current user is a superuser. (CVE-2017-7486)

By itself, this patch will only fix the behavior in newly initdb’d databases. If you wish to apply this change in an existing database, follow the corrected procedure shown in the changelog entry for CVE-2017-7547, in Section E.14.

• Prevent exposure of statistical information via leaky operators (Peter Eisentraut)

Some selectivity estimation functions in the planner will apply user-defined operators to values obtained from pg_statistic, such as most common values and histogram entries. This occurs before table permissions are checked, so a nefarious user could exploit the behavior to obtain these values for table columns he does not have permission to read. To fix, fall back to a default estimate if the operator’s implementation function is not certified leak-proof and the calling user does not have permission to read the table column whose statistics are needed. At least one of these criteria is satisfied in most cases in practice. (CVE-2017-7484)

• Restore libpq’s recognition of the PGREQUIRESSL environment variable (Daniel Gustafsson)

Processing of this environment variable was unintentionally dropped in PostgreSQL 9.3, but its documentation remained. This creates a security hazard, since users might be relying on the environment variable to force SSL-encrypted connections, but that would no longer be guaranteed.
Appendix E. Release Notes

Restore handling of the variable, but give it lower priority than PGSSLMODE, to avoid breaking configurations that work correctly with post-9.3 code. (CVE-2017-7485)

- Fix possibly-invalid initial snapshot during logical decoding (Petr Jelinek, Andres Freund)
  The initial snapshot created for a logical decoding replication slot was potentially incorrect. This could cause third-party tools that use logical decoding to copy incomplete/inconsistent initial data. This was more likely to happen if the source server was busy at the time of slot creation, or if another logical slot already existed.

  If you are using a replication tool that depends on logical decoding, and it should have copied a nonempty data set at the start of replication, it is advisable to recreate the replica after installing this update, or to verify its contents against the source server.

- Fix possible corruption of “init forks” of unlogged indexes (Robert Haas, Michael Paquier)
  This could result in an unlogged index being set to an invalid state after a crash and restart. Such a problem would persist until the index was dropped and rebuilt.

- Fix incorrect reconstruction of pg_subtrans entries when a standby server replays a prepared but uncommitted two-phase transaction (Tom Lane)
  In most cases this turned out to have no visible ill effects, but in corner cases it could result in circular references in pg_subtrans, potentially causing infinite loops in queries that examine rows modified by the two-phase transaction.

- Avoid possible crash in walsender due to failure to initialize a string buffer (Stas Kelvich, Fujii Masao)

- Fix possible crash when rescanning a nearest-neighbor index-only scan on a GiST index (Tom Lane)

- Prevent delays in postmaster’s launching of multiple parallel worker processes (Tom Lane)
  There could be a significant delay (up to tens of seconds) before satisfying a query’s request for more than one worker process, or when multiple queries requested workers simultaneously. On most platforms this required unlucky timing, but on some it was the typical case.

- Fix postmaster’s handling of fork() failure for a background worker process (Tom Lane)
  Previously, the postmaster updated portions of its state as though the process had been launched successfully, resulting in subsequent confusion.

- Fix possible “no relation entry for relid 0” error when planning nested set operations (Tom Lane)

- Fix assorted minor issues in planning of parallel queries (Robert Haas)

- Avoid applying “physical targetlist” optimization to custom scans (Dmitry Ivanov, Tom Lane)
  This optimization supposed that retrieving all columns of a tuple is inexpensive, which is true for ordinary Postgres tuples; but it might not be the case for a custom scan provider.

- Use the correct sub-expression when applying a FOR ALL row-level-security policy (Stephen Frost)
  In some cases the WITH CHECK restriction would be applied when the USING restriction is more appropriate.

- Ensure parsing of queries in extension scripts sees the results of immediately-preceding DDL (Julien Rouhaud, Tom Lane)
  Due to lack of a cache flush step between commands in an extension script file, non-utility queries might not see the effects of an immediately preceding catalog change, such as ALTER TABLE ... RENAME.
• Skip tablespace privilege checks when `ALTER TABLE ... ALTER COLUMN TYPE` rebuilds an existing index (Noah Misch)

The command failed if the calling user did not currently have `CREATE` privilege for the tablespace containing the index. That behavior seems unhelpful, so skip the check, allowing the index to be rebuilt where it is.

• Fix `ALTER TABLE ... VALIDATE CONSTRAINT` to not recurse to child tables when the constraint is marked `NO INHERIT` (Amit Langote)

This fix prevents unwanted “constraint does not exist” failures when no matching constraint is present in the child tables.

• Avoid dangling pointer in `COPY ... TO` when row-level security is active for the source table (Tom Lane)

Usually this had no ill effects, but sometimes it would cause unexpected errors or crashes.

• Avoid accessing an already-closed relcache entry in `CLUSTER` and `VACUUM FULL` (Tom Lane)

With some bad luck, this could lead to indexes on the target relation getting rebuilt with the wrong persistence setting.

• Fix `VACUUM` to account properly for pages that could not be scanned due to conflicting page pins (Andrew Gierth)

This tended to lead to underestimation of the number of tuples in the table. In the worst case of a small heavily-contended table, `VACUUM` could incorrectly report that the table contained no tuples, leading to very bad planning choices.

• Ensure that bulk-tuple-transfer loops within a hash join are interruptible by query cancel requests (Tom Lane, Thomas Munro)

• Fix incorrect support for certain `box` operators in SP-GiST (Nikita Glukhov)

SP-GiST index scans using the operators `&< &> &<|` and `&>|` would yield incorrect answers.

• Fix integer-overflow problems in `interval` comparison (Kyotaro Horiguchi, Tom Lane)

The comparison operators for type `interval` could yield wrong answers for intervals larger than about 296000 years. Indexes on columns containing such large values should be reindexed, since they may be corrupt.

• Fix `cursor_to_xml()` to produce valid output with `tableforest = false` (Thomas Munro, Peter Eisentraut)

Previously it failed to produce a wrapping `<table>` element.

• Fix roundoff problems in `float8_timestamptz()` and `make_interval()` (Tom Lane)

These functions truncated, rather than rounded, when converting a floating-point value to integer microseconds; that could cause unexpectedly off-by-one results.

• Fix `pg_get_object_address()` to handle members of operator families correctly (Álvaro Herrera)

• Fix cancelling of `pg_stop_backup()` when attempting to stop a non-exclusive backup (Michael Paquier, David Steele)

If `pg_stop_backup()` was cancelled while waiting for a non-exclusive backup to end, related state was left inconsistent; a new exclusive backup could not be started, and there were other minor problems.

• Improve performance of `pg_timezone_names` view (Tom Lane, David Rowley)

• Reduce memory management overhead for contexts containing many large blocks (Tom Lane)
• Fix sloppy handling of corner-case errors from `lseek()` and `close()` (Tom Lane)
  Neither of these system calls are likely to fail in typical situations, but if they did, `fd.c` could get quite confused.

• Fix incorrect check for whether postmaster is running as a Windows service (Michael Paquier)
  This could result in attempting to write to the event log when that’s not accessible, so that no logging happens at all.

• Fix `ecpg` to support `COMMIT PREPARED` and `ROLLBACK PREPARED` (Masahiko Sawada)

• Fix a double-free error when processing dollar-quoted string literals in `ecpg` (Michael Meskes)

• Fix `pgbench` to handle the combination of `--connect` and `--rate` options correctly (Fabien Coelho)

• Fix `pgbench` to honor the long-form option spelling `--builtin`, as per its documentation (Tom Lane)

• Fix `pg_dump/pg_restore` to correctly handle privileges for the `public` schema when using `--clean` option (Stephen Frost)
  Other schemas start out with no privileges granted, but `public` does not; this requires special-case treatment when it is dropped and restored due to the `--clean` option.

• In `pg_dump`, fix incorrect schema and owner marking for comments and security labels of some types of database objects (Giuseppe Broccolo, Tom Lane)
  In simple cases this caused no ill effects; but for example, a schema-selective restore might omit comments it should include, because they were not marked as belonging to the schema of their associated object.

• Fix typo in `pg_dump`’s query for initial privileges of a procedural language (Peter Eisentraut)
  This resulted in `pg_dump` always believing that the language had no initial privileges. Since that’s true for most procedural languages, ill effects from this bug are probably rare.

• Avoid emitting an invalid list file in `pg_restore -l` when SQL object names contain newlines (Tom Lane)
  Replace newlines by spaces, which is sufficient to make the output valid for `pg_restore -L`’s purposes.

• Fix `pg_upgrade` to transfer comments and security labels attached to “large objects” (blobs) (Stephen Frost)
  Previously, blobs were correctly transferred to the new database, but any comments or security labels attached to them were lost.

• Improve error handling in `contrib/adminpack`'s `pg_file_write()` function (Noah Misch)
  Notably, it failed to detect errors reported by `fclose()`.

• In `contrib/dblink`, avoid leaking the previous unnamed connection when establishing a new unnamed connection (Joe Conway)

• Fix `contrib/pg_trgm`'s extraction of trigrams from regular expressions (Tom Lane)
  In some cases it would produce a broken data structure that could never match anything, leading to GIN or GiST index scans that use a trigram index not finding any matches to the regular expression.

• In `contrib/postgres_fdw`, allow join conditions that contain shippable extension-provided functions to be pushed to the remote server (David Rowley, Ashutosh Bapat)

• Support Tcl 8.6 in MSVC builds (Álvaro Herrera)
- Sync our copy of the timezone library with IANA release tzcode2017b (Tom Lane)
  This fixes a bug affecting some DST transitions in January 2038.
- Update time zone data files to tzdata release 2017b for DST law changes in Chile, Haiti, and Mongolia, plus historical corrections for Ecuador, Kazakhstan, Liberia, and Spain. Switch to numeric abbreviations for numerous time zones in South America, the Pacific and Indian oceans, and some Asian and Middle Eastern countries.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the `pg_timezone_names` view nor used for output.

- Use correct daylight-savings rules for POSIX-style time zone names in MSVC builds (David Rowley)

The Microsoft MSVC build scripts neglected to install the `posixrules` file in the timezone directory tree. This resulted in the timezone code falling back to its built-in rule about what DST behavior to assume for a POSIX-style time zone name. For historical reasons that still corresponds to the DST rules the USA was using before 2007 (i.e., change on first Sunday in April and last Sunday in October). With this fix, a POSIX-style zone name will use the current and historical DST transition dates of the `US/Eastern` zone. If you don’t want that, remove the `posixrules` file, or replace it with a copy of some other zone file (see Section 8.5.3). Note that due to caching, you may need to restart the server to get such changes to take effect.

### E.16. Release 9.6.2

**Release date:** 2017-02-09

This release contains a variety of fixes from 9.6.1. For information about new features in the 9.6 major release, see Section E.18.

#### E.16.1. Migration to Version 9.6.2

A dump/restore is not required for those running 9.6.X.

However, if your installation has been affected by the bug described in the first changelog entry below, then after updating you may need to take action to repair corrupted indexes.

Also, if you are upgrading from a version earlier than 9.6.1, see Section E.17.

#### E.16.2. Changes

- Fix a race condition that could cause indexes built with `CREATE INDEX CONCURRENTLY` to be corrupt (Pavan Deolasee, Tom Lane)

  If `CREATE INDEX CONCURRENTLY` was used to build an index that depends on a column not previously indexed, then rows updated by transactions that ran concurrently with the `CREATE INDEX`
command could have received incorrect index entries. If you suspect this may have happened, the most reliable solution is to rebuild affected indexes after installing this update.

- Ensure that the special snapshot used for catalog scans is not invalidated by premature data pruning (Tom Lane)
  Backends failed to account for this snapshot when advertising their oldest xmin, potentially allowing concurrent vacuuming operations to remove data that was still needed. This led to transient failures along the lines of “cache lookup failed for relation 1255”.

- Fix incorrect WAL logging for BRIN indexes (Kuntal Ghosh)
  The WAL record emitted for a BRIN “revmap” page when moving an index tuple to a different page was incorrect. Replay would make the related portion of the index useless, forcing it to be recomputed.

- Unconditionally WAL-log creation of the “init fork” for an unlogged table (Michael Paquier)
  Previously, this was skipped when wal_level = minimal, but actually it’s necessary even in that case to ensure that the unlogged table is properly reset to empty after a crash.

- If the stats collector dies during hot standby, restart it (Takayuki Tsunakawa)

- Ensure that hot standby feedback works correctly when it’s enabled at standby server start (Ants Aasma, Craig Ringer)

- Check for interrupts while hot standby is waiting for a conflicting query (Simon Riggs)

- Avoid constantly respawning the autovacuum launcher in a corner case (Amit Khandekar)
  This fix avoids problems when autovacuum is nominally off and there are some tables that require freezing, but all such tables are already being processed by autovacuum workers.

- Disallow setting the num_sync field to zero in synchronous_standby_names (Fujii Masao)
  The correct way to disable synchronous standby is to set the whole value to an empty string.

- Don’t count background worker processes against a user’s connection limit (David Rowley)

- Fix check for when an extension member object can be dropped (Tom Lane)
  Extension upgrade scripts should be able to drop member objects, but this was disallowed for serial-column sequences, and possibly other cases.

- Fix tracking of initial privileges for extension member objects so that it works correctly with ALTER EXTENSION ... ADD/DROP (Stephen Frost)
  An object’s current privileges at the time it is added to the extension will now be considered its default privileges; only later changes in its privileges will be dumped by subsequent pg_dump runs.

- Make sure ALTER TABLE preserves index tablespace assignments when rebuilding indexes (Tom Lane, Michael Paquier)
  Previously, non-default settings of default_tablespace could result in broken indexes.

- Fix incorrect updating of trigger function properties when changing a foreign-key constraint’s deferrability properties with ALTER TABLE ... ALTER CONSTRAINT (Tom Lane)
  This led to odd failures during subsequent exercise of the foreign key, as the triggers were fired at the wrong times.

- Prevent dropping a foreign-key constraint if there are pending trigger events for the referenced relation (Tom Lane)
  This avoids “could not find trigger NNN” or “relation NNN has no triggers” errors.
• Fix `ALTER TABLE ... SET DATA TYPE ... USING` when child table has different column ordering than the parent (Álvaro Herrera)

Failure to adjust the column numbering in the `USING` expression led to errors, typically “attribute \(N\) has wrong type”.

• Fix processing of OID column when a table with OIDs is associated to a parent with OIDs via `ALTER TABLE ... INHERIT` (Amit Langote)

The OID column should be treated the same as regular user columns in this case, but it wasn’t, leading to odd behavior in later inheritance changes.

• Ensure that `CREATE TABLE ... LIKE ... WITH OIDS` creates a table with OIDs, whether or not the `LIKE`-referenced table(s) have OIDs (Tom Lane)

• Fix `CREATE OR REPLACE VIEW` to update the view query before attempting to apply the new view options (Dean Rasheed)

Previously the command would fail if the new options were inconsistent with the old view definition.

• Report correct object identity during `ALTER TEXT SEARCH CONFIGURATION` (Artur Zakirov)

The wrong catalog OID was reported to extensions such as logical decoding.

• Fix commit timestamp mechanism to not fail when queried about the special XIDs `FrozenTransactionId` and `BootstrapTransactionId` (Craig Ringer)

• Fix incorrect use of view reloptions as regular table reloptions (Tom Lane)

The symptom was spurious “ON CONFLICT is not supported on table ... used as a catalog table” errors when the target of `INSERT ... ON CONFLICT` is a view with cascade option.

• Fix incorrect “target lists can have at most \(N\) entries” complaint when using `ON CONFLICT` with wide tables (Tom Lane)

• Fix spurious “query provides a value for a dropped column” errors during `INSERT` or `UPDATE` on a table with a dropped column (Tom Lane)

• Prevent multicolumn expansion of `foo.*` in an `UPDATE` source expression (Tom Lane)

This led to “UPDATE target count mismatch --- internal error”. Now the syntax is understood as a whole-row variable, as it would be in other contexts.

• Ensure that column typmods are determined accurately for multi-row `VALUES` constructs (Tom Lane)

This fixes problems occurring when the first value in a column has a determinable typmod (e.g., length for a `varchar` value) but later values don’t share the same limit.

• Throw error for an unfinished Unicode surrogate pair at the end of a Unicode string (Tom Lane)

Normally, a Unicode surrogate leading character must be followed by a Unicode surrogate trailing character, but the check for this was missed if the leading character was the last character in a Unicode string literal (`U&'...'`) or Unicode identifier (`U&"..."`).

• Fix execution of `DISTINCT` and ordered aggregates when multiple such aggregates are able to share the same transition state (Heikki Linnakangas)

• Fix implementation of phrase search operators in `tsquery` (Tom Lane)

Remove incorrect, and inconsistently-applied, rewrite rules that tried to transform away AND/OR/NOT operators appearing below a `PHRASE` operator; instead upgrade the execution engine to handle such cases correctly. This fixes assorted strange behavior and possible crashes for text search queries containing such combinations. Also fix nested `PHRASE` operators to work
Appendix E. Release Notes

sanely in combinations other than simple left-deep trees, correct the behavior when removing stopwords from a phrase search clause, and make sure that index searches behave consistently with simple sequential-scan application of such queries.

- Ensure that a purely negative text search query, such as !foo, matches empty tsvectors (Tom Dunstan)

  Such matches were found by GIN index searches, but not by sequential scans or GiST index searches.

- Prevent crash when ts_rewrite() replaces a non-top-level subtree with an empty query (Artur Zakirov)

- Fix performance problems in ts_rewrite() (Tom Lane)

- Fix ts_rewrite()’s handling of nested NOT operators (Tom Lane)

- Improve speed of user-defined aggregates that use array_append() as transition function (Tom Lane)

- Fix array_fill() to handle empty arrays properly (Tom Lane)

- Fix possible crash in array_position() or array_positions() when processing arrays of records (Junseok Yang)

- Fix one-byte buffer overrun in quote_literal_cstr() (Heikki Linnakangas)

  The overrun occurred only if the input consisted entirely of single quotes and/or backslashes.

- Prevent multiple calls of pg_start_backup() and pg_stop_backup() from running concurrently (Michael Paquier)

  This avoids an assertion failure, and possibly worse things, if someone tries to run these functions in parallel.

- Disable transform that attempted to remove no-op AT TIME ZONE conversions (Tom Lane)

  This resulted in wrong answers when the simplified expression was used in an index condition.

- Avoid discarding interval-to-interval casts that aren’t really no-ops (Tom Lane)

  In some cases, a cast that should result in zeroing out low-order interval fields was mistakenly deemed to be a no-op and discarded. An example is that casting from INTERVAL MONTH to INTERVAL YEAR failed to clear the months field.

- Fix crash if the number of workers available to a parallel query decreases during a rescan (Andreas Seltenreich)

- Fix bugs in transmitting GUC parameter values to parallel workers (Michael Paquier, Tom Lane)

- Allow statements prepared with PREPARE to be given parallel plans (Amit Kapila, Tobias Bussmann)

- Fix incorrect generation of parallel plans for semi-joins (Tom Lane)

- Fix planner’s cardinality estimates for parallel joins (Robert Haas)

  Ensure that these estimates reflect the number of rows predicted to be seen by each worker, rather than the total.

- Fix planner to avoid trying to parallelize plan nodes containing initplans or subplans (Tom Lane, Amit Kapila)

- Ensure that cached plans are invalidated by changes in foreign-table options (Amit Langote, Etsuro Fujita, Ashutosh Bapat)

- Fix the plan generated for sorted partial aggregation with a constant GROUP BY clause (Tom Lane)
Appendix E. Release Notes

- Fix "could not find plan for CTE" planner error when dealing with a `UNION ALL` containing CTE references (Tom Lane)
- Fix mishandling of initplans when forcibly adding a Material node to a subplan (Tom Lane)
  The typical consequence of this mistake was a "plan should not reference subplan’s variable" error.
- Fix foreign-key-based join selectivity estimation for semi-joins and anti-joins, as well as inheritance cases (Tom Lane)
  The new code for taking the existence of a foreign key relationship into account did the wrong thing in these cases, making the estimates worse not better than the pre-9.6 code.
- Fix `pg_dump` to emit the data of a sequence that is marked as an extension configuration table (Michael Paquier)
- Fix mishandling of `ALTER DEFAULT PRIVILEGES ... REVOKE` in `pg_dump` (Stephen Frost)
  `pg_dump` missed issuing the required `REVOKE` commands in cases where `ALTER DEFAULT PRIVILEGES` had been used to reduce privileges to less than they would normally be.
- Fix `pg_dump` to dump user-defined casts and transforms that use built-in functions (Stephen Frost)
- Fix `pg_restore` with `--create --if-exists` to behave more sanely if an archive contains unrecognized `DROP` commands (Tom Lane)
  This doesn’t fix any live bug, but it may improve the behavior in future if `pg_restore` is used with an archive generated by a later `pg_dump` version.
- Fix `pg_basebackup`’s rate limiting in the presence of slow I/O (Antonin Houska)
  If disk I/O was transiently much slower than the specified rate limit, the calculation overflowed, effectively disabling the rate limit for the rest of the run.
- Fix `pg_basebackup`’s handling of symlinked `pg_stat_tmp` and `pg_replslot` subdirectories (Magnus Hagander, Michael Paquier)
- Fix possible `pg_basebackup` failure on standby server when including WAL files (Amit Kapila, Robert Haas)
- Improve `initdb` to insert the correct platform-specific default values for the `xxx_flush_after` parameters into `postgresql.conf` (Fabien Coelho, Tom Lane)
  This is a cleaner way of documenting the default values than was used previously.
- Fix possible mishandling of expanded arrays in domain check constraints and `CASE` execution (Tom Lane)
  It was possible for a PL/pgSQL function invoked in these contexts to modify or even delete an array value that needs to be preserved for additional operations.
- Fix nested uses of PL/pgSQL functions in contexts such as domain check constraints evaluated during assignment to a PL/pgSQL variable (Tom Lane)
- Ensure that the Python exception objects we create for PL/Python are properly reference-counted (Rafa de la Torre, Tom Lane)
  This avoids failures if the objects are used after a Python garbage collection cycle has occurred.
- Fix PL/Tcl to support triggers on tables that have `.tupno` as a column name (Tom Lane)
  This matches the (previously undocumented) behavior of PL/Tcl’s `spi_exec` and `spi_execp` commands, namely that a magic `.tupno` column is inserted only if there isn’t a real column named that.
- Allow DOS-style line endings in `~/.pgpass` files, even on Unix (Vik Fearing)
This change simplifies use of the same password file across Unix and Windows machines.

- Fix one-byte buffer overrun if ecpg is given a file name that ends with a dot (Takayuki Tsunakawa)
- Fix incorrect error reporting for duplicate data in psqI’s \crosstabview (Tom Lane)
- psql sometimes quoted the wrong row and/or column values when complaining about multiple entries for the same crosstab cell.
- Fix psql’s tab completion for ALTER DEFAULT PRIVILEGES (Gilles Darold, Stephen Frost)
- Fix psql’s tab completion for ALTER TABLE t ALTER c DROP . . . (Kyotaro Horiguchi)
- In psql, treat an empty or all-blank setting of the PAGER environment variable as meaning “no pager” (Tom Lane)
  Previously, such a setting caused output intended for the pager to vanish entirely.
- Improve contrib/dblink’s reporting of low-level libpq errors, such as out-of-memory (Joe Con- way)
- Teach contrib/dblink to ignore irrelevant server options when it uses a contrib/postgres_fdw foreign server as the source of connection options (Corey Huinker)
  Previously, if the foreign server object had options that were not also libpq connection options, an error occurred.
- Fix portability problems in contrib/pageinspect’s functions for GIN indexes (Peter Eisentraut, Tom Lane)
- Fix possible miss of socket read events while waiting on Windows (Amit Kapila)
  This error was harmless for most uses, but it is known to cause hangs when trying to use the pldebugger extension.
- On Windows, ensure that environment variable changes are propagated to DLLs built with debug options (Christian Ullrich)
- Sync our copy of the timezone library with IANA release tzcode2016j (Tom Lane)
  This fixes various issues, most notably that timezone data installation failed if the target directory didn’t support hard links.
- Update time zone data files to tzdata release 2016j for DST law changes in northern Cyprus (adding a new zone Asia/Famagusta), Russia (adding a new zone Europe/Saratov), Tonga, and Antarc- tica/Casey. Historical corrections for Italy, Kazakhstan, Malta, and Palestine. Switch to preferring numeric zone abbreviations for Tonga.

### E.17. Release 9.6.1

**Release date:** 2016-10-27

This release contains a variety of fixes from 9.6.0. For information about new features in the 9.6 major release, see Section E.18.

#### E.17.1. Migration to Version 9.6.1

A dump/restore is not required for those running 9.6.X.
However, if your installation has been affected by the bugs described in the first two changelog entries below, then after updating you may need to take action to repair corrupted free space maps and/or visibility maps.

**E.17.2. Changes**

- **Fix WAL-logging of truncation of relation free space maps and visibility maps (Pavan Deolasee, Heikki Linnakangas)**
  It was possible for these files to not be correctly restored during crash recovery, or to be written incorrectly on a standby server. Bogus entries in a free space map could lead to attempts to access pages that have been truncated away from the relation itself, typically producing errors like “could not read block XXX: read only 0 of 8192 bytes”. Checksum failures in the visibility map are also possible, if checksumming is enabled.
  Procedures for determining whether there is a problem and repairing it if so are discussed at https://wiki.postgresql.org/wiki/Free_Space_Map_Problems.
- **Fix possible data corruption when pg_upgrade rewrites a relation visibility map into 9.6 format (Tom Lane)**
  On big-endian machines, bytes of the new visibility map were written in the wrong order, leading to a completely incorrect map. On Windows, the old map was read using text mode, leading to incorrect results if the map happened to contain consecutive bytes that matched a carriage return/line feed sequence. The latter error would almost always lead to a pg_upgrade failure due to the map file appearing to be the wrong length.
  If you are using a big-endian machine (many non-Intel architectures are big-endian) and have used pg_upgrade to upgrade from a pre-9.6 release, you should assume that all visibility maps are incorrect and need to be regenerated. It is sufficient to truncate each relation’s visibility map with contrib/pg_visibility’s pg_truncate_visibility_map() function. For more information see https://wiki.postgresql.org/wiki/Visibility_Map_Problems.
- **Don’t throw serialization errors for self-conflicting insertions in INSERT ... ON CONFLICT (Thomas Munro, Peter Geoghegan)**
- **Fix use-after-free hazard in execution of aggregate functions using DISTINCT (Peter Geoghegan)**
  This could lead to a crash or incorrect query results.
- **Fix incorrect handling of polymorphic aggregates used as window functions (Tom Lane)**
  The aggregate’s transition function was told that its first argument and result were of the aggregate’s output type, rather than the state type. This led to errors or crashes with polymorphic transition functions.
- **Fix COPY with a column name list from a table that has row-level security enabled (Adam Brightwell)**
- **Fix EXPLAIN to emit valid XML when track_io_timing is on (Markus Winand)**
  Previously the XML output-format option produced syntactically invalid tags such as `<I/O-Read-Time>`. That is now rendered as `<I-O-Read-Time>`.
- **Fix statistics update for TRUNCATE in a prepared transaction (Stas Kelvich)**
- **Fix bugs in merging inherited CHECK constraints while creating or altering a table (Tom Lane, Amit Langote)**
Appendix E. Release Notes

Allow identical \texttt{CHECK} constraints to be added to a parent and child table in either order. Prevent merging of a valid constraint from the parent table with a \texttt{NOT VALID} constraint on the child. Likewise, prevent merging of a \texttt{NO INHERIT} child constraint with an inherited constraint.

- Show a sensible value in \texttt{pg_settings.unit} for \texttt{min_wal_size} and \texttt{max_wal_size} (Tom Lane)

- Fix replacement of array elements in \texttt{jsonb_set()} (Tom Lane)
  If the target is an existing JSON array element, it got deleted instead of being replaced with a new value.

- Avoid very-low-probability data corruption due to testing tuple visibility without holding buffer lock (Thomas Munro, Peter Geoghegan, Tom Lane)

- Preserve commit timestamps across server restart (Julien Rouhaud, Craig Ringer)
  With \texttt{track_commit_timestamp} turned on, old commit timestamps became inaccessible after a clean server restart.

- Fix logical WAL decoding to work properly when a subtransaction’s WAL output is large enough to spill to disk (Andres Freund)

- Fix dangling-pointer problem in logical WAL decoding (Stas Kelvich)

- Round shared-memory allocation request to a multiple of the actual huge page size when attempting to use huge pages on Linux (Tom Lane)
  This avoids possible failures during \texttt{munmap()} on systems with atypical default huge page sizes. Except in crash-recovery cases, there were no ill effects other than a log message.

- Don’t try to share SSL contexts across multiple connections in \texttt{libpq} (Heikki Linnakangas)
  This led to assorted corner-case bugs, particularly when trying to use different SSL parameters for different connections.

- Avoid corner-case memory leak in \texttt{libpq} (Tom Lane)
  The reported problem involved leaking an error report during \texttt{PQreset()}, but there might be related cases.

- In \texttt{pg_upgrade}, check library loadability in name order (Tom Lane)
  This is a workaround to deal with cross-extension dependencies from language transform modules to their base language and data type modules.

- Fix \texttt{pg_upgrade} to work correctly for extensions containing index access methods (Tom Lane)
  To allow this, the server has been extended to support \texttt{ALTER EXTENSION ADD/DROP ACCESS METHOD}. That functionality should have been included in the original patch to support dynamic creation of access methods, but it was overlooked.

- Improve error reporting in \texttt{pg_upgrade}’s file copying/linking/rewriting steps (Tom Lane, Álvaro Herrera)

- Fix \texttt{pg_dump} to work against pre-7.4 servers (Amit Langote, Tom Lane)

- Disallow specifying both \texttt{--source-server} and \texttt{--source-target} options to \texttt{pg_rewind} (Michael Banck)

- Make \texttt{pg_rewind} turn off \texttt{synchronous_commit} in its session on the source server (Michael Banck, Michael Paquier)
  This allows \texttt{pg_rewind} to work even when the source server is using synchronous replication that is not working for some reason.
Appendix E. Release Notes

• In pg_xlogdump, retry opening new WAL segments when using --follow option (Magnus Hagander)
  This allows for a possible delay in the server’s creation of the next segment.
• Fix contrib/pg_visibility to report the correct TID for a corrupt tuple that has been the subject of a rolled-back update (Tom Lane)
• Fix makefile dependencies so that parallel make of PL/Python by itself will succeed reliably (Pavel Raiskup)
• Update time zone data files to tzdata release 2016h for DST law changes in Palestine and Turkey, plus historical corrections for Turkey and some regions of Russia. Switch to numeric abbreviations for some time zones in Antarctica, the former Soviet Union, and Sri Lanka.

The IANA time zone database previously provided textual abbreviations for all time zones, sometimes making up abbreviations that have little or no currency among the local population. They are in process of reversing that policy in favor of using numeric UTC offsets in zones where there is no evidence of real-world use of an English abbreviation. At least for the time being, PostgreSQL will continue to accept such removed abbreviations for timestamp input. But they will not be shown in the pg_timezone_names view nor used for output.

In this update, AMT is no longer shown as being in use to mean Armenia Time. Therefore, we have changed the Default abbreviation set to interpret it as Amazon Time, thus UTC-4 not UTC+4.

E.18. Release 9.6
Release date: 2016-09-29

E.18.1. Overview

Major enhancements in PostgreSQL 9.6 include:

• Parallel execution of sequential scans, joins and aggregates
• Avoid scanning pages unnecessarily during vacuum freeze operations
• Synchronous replication now allows multiple standby servers for increased reliability
• Full-text search can now search for phrases (multiple adjacent words)
• postgres_fdw now supports remote joins, sorts, UPDATES, and DELETES
• Substantial performance improvements, especially in the area of scalability on multi-CPU-socket servers

The above items are explained in more detail in the sections below.

E.18.2. Migration to Version 9.6

A dump/restore using pg_dumpall, or use of pg_upgrade, is required for those wishing to migrate data from any previous release.

Version 9.6 contains a number of changes that may affect compatibility with previous releases. Observe the following incompatibilities:
Appendix E. Release Notes

- Improve the `pg_stat_activity` view’s information about what a process is waiting for (Amit Kapila, Ildus Kurbangaliev)

  Historically a process has only been shown as waiting if it was waiting for a heavyweight lock. Now waits for lightweight locks and buffer pins are also shown in `pg_stat_activity`. Also, the type of lock being waited for is now visible. These changes replace the `waiting` column with `wait_event_type` and `wait_event`.

- In `to_char()`, do not count a minus sign (when needed) as part of the field width for time-related fields (Bruce Momjian)

  For example, `to_char('-4 years'::interval, 'YY')` now returns `-04`, rather than `-4`.

- Make `extract()` behave more reasonably with infinite inputs (Vitaly Burovoy)

  Historically the `extract()` function just returned zero given an infinite timestamp, regardless of the given field name. Make it return `infinity` or `-infinity` as appropriate when the requested field is one that is monotonically increasing (e.g., `year`, `epoch`), or `NULL` when it is not (e.g., `day`, `hour`). Also, throw the expected error for bad field names.

- Remove PL/pgSQL’s “feature” that suppressed the innermost line of `CONTEXT` for messages emitted by `RAISE` commands (Pavel Stehule)

  This ancient backwards-compatibility hack was agreed to have outlived its usefulness.

- Fix the default text search parser to allow leading digits in `email` and `host` tokens (Artur Zakirov)

  In most cases this will result in few changes in the parsing of text. But if you have data where such addresses occur frequently, it may be worth rebuilding dependent `tsvector` columns and indexes so that addresses of this form will be found properly by text searches.

- Extend `contrib/unaccent`’s standard `unaccent.rules` file to handle all diacritics known to Unicode, and to expand ligatures correctly (Thomas Munro, Léonard Benedetti)

  The previous version neglected to convert some less-common letters with diacritic marks. Also, ligatures are now expanded into separate letters. Installations that use this rules file may wish to rebuild `tsvector` columns and indexes that depend on the result.

- Remove the long-deprecated `CREATEUSER/NOCREATEUSER` options from `CREATE ROLE` and allied commands (Tom Lane)

  `CREATEUSER` actually meant `SUPERUSER`, for ancient backwards-compatibility reasons. This has been a constant source of confusion for people who (reasonably) expect it to mean `CREATEROLE`. It has been deprecated for ten years now, so fix the problem by removing it.

- Treat role names beginning with `pg_` as reserved (Stephen Frost)

  User creation of such role names is now disallowed. This prevents conflicts with built-in roles created by `initdb`.

- Change a column name in the `information_schema.routines` view from `result_cast_character_set_name` to `result_cast_char_set_name` (Clément Prévost)

  The SQL:2011 standard specifies the longer name, but that appears to be a mistake, because adjacent column names use the shorter style, as do other `information_schema` views.

- `psql`’s `-c` option no longer implies `--no-psqlrc` (Pavel Stehule, Catalin Iacob)

  Write `--no-psqlrc` (or its abbreviation `-X`) explicitly to obtain the old behavior. Scripts so modified will still work with old versions of `psql`.

- Improve `pg_restore`’s `-t` option to match all types of relations, not only plain tables (Craig Ringer)
Appendix E. Release Notes

- Change the display format used for NextXID in pg_controldata and related places (Joe Conway, Bruce Momjian)
  
  Display epoch-and-transaction-ID values in the format number:number. The previous format number/number was confusingly similar to that used for LSNs.

- Update extension functions to be marked parallel-safe where appropriate (Andreas Karlsson)
  
  Many of the standard extensions have been updated to allow their functions to be executed within parallel query worker processes. These changes will not take effect in databases pg_upgrade’d from prior versions unless you apply ALTER EXTENSION UPDATE to each such extension (in each database of a cluster).

E.18.3. Changes

Below you will find a detailed account of the changes between PostgreSQL 9.6 and the previous major release.

E.18.3.1. Server

E.18.3.1.1. Parallel Queries

- Parallel queries (Robert Haas, Amit Kapila, David Rowley, many others)
  
  With 9.6, PostgreSQL introduces initial support for parallel execution of large queries. Only strictly read-only queries where the driving table is accessed via a sequential scan can be parallelized. Hash joins and nested loops can be performed in parallel, as can aggregation (for supported aggregates). Much remains to be done, but this is already a useful set of features.

  Parallel query execution is not (yet) enabled by default. To allow it, set the new configuration parameter max_parallel_workers_per_gather to a value larger than zero. Additional control over use of parallelism is available through other new configuration parameters force_parallel_mode, parallel_setup_cost, parallel_tuple_cost, and min_parallel_relation_size.

- Provide infrastructure for marking the parallel-safety status of functions (Robert Haas, Amit Kapila)

E.18.3.1.2. Indexes

- Allow GIN index builds to make effective use of maintenance_work_mem settings larger than 1 GB (Robert Abraham, Teodor Sigaev)

- Add pages deleted from a GIN index’s pending list to the free space map immediately (Jeff Janes, Teodor Sigaev)

  This reduces bloat if the table is not vacuumed often.

- Add gin_clean_pending_list() function to allow manual invocation of pending-list cleanup for a GIN index (Jeff Janes)

  Formerly, such cleanup happened only as a byproduct of vacuuming or analyzing the parent table.

- Improve handling of dead index tuples in GiST indexes (Anastasia Lubennikova)
Dead index tuples are now marked as such when an index scan notices that the corresponding heap tuple is dead. When inserting tuples, marked-dead tuples will be removed if needed to make space on the page.

- Add an SP-GiST operator class for type box (Alexander Lebedev)

### E.18.3.1.3. Sorting

- Improve sorting performance by using quicksort, not replacement selection sort, when performing external sort steps (Peter Geoghegan)

  The new approach makes better use of the CPU cache for typical cache sizes and data volumes. Where necessary, the behavior can be adjusted via the new configuration parameter replacement_sort_tuples.

- Speed up text sorts where the same string occurs multiple times (Peter Geoghegan)

- Speed up sorting of uuid, bytea, and char(n) fields by using “abbreviated” keys (Peter Geoghegan)

  Support for abbreviated keys has also been added to the non-default operator classes text_pattern_ops, varchar_pattern_ops, and bpchar_pattern_ops. Processing of ordered-set aggregates can also now exploit abbreviated keys.

- Speed up `CREATE INDEX CONCURRENTLY` by treating TIDs as 64-bit integers during sorting (Peter Geoghegan)

### E.18.3.1.4. Locking

- Reduce contention for the ProcArrayLock (Amit Kapila, Robert Haas)

- Improve performance by moving buffer content locks into the buffer descriptors (Andres Freund, Simon Riggs)

- Replace shared-buffer header spinlocks with atomic operations to improve scalability (Alexander Korotkov, Andres Freund)

- Use atomic operations, rather than a spinlock, to protect an LWLock’s wait queue (Andres Freund)

- Partition the shared hash table freelist to reduce contention on multi-CPU-socket servers (Aleksander Alekseev)

- Reduce interlocking on standby servers during the replay of btree index vacuuming operations (Simon Riggs)

  This change avoids substantial replication delays that sometimes occurred while replaying such operations.

### E.18.3.1.5. Optimizer Statistics

- Improve `ANALYZE`’s estimates for columns with many nulls (Tomas Vondra, Alex Shulgin)

  Previously `ANALYZE` tended to underestimate the number of non-NULL distinct values in a column with many NULLs, and was also inaccurate in computing the most-common values.

- Improve planner’s estimate of the number of distinct values in a query result (Tomas Vondra)

- Use foreign key relationships to infer selectivity for join predicates (Tomas Vondra, David Rowley)
Appendix E. Release Notes

If a table \( t \) has a foreign key restriction, say \((a,b) \text{ REFERENCES } r \ (x,y)\), then a \texttt{WHERE} condition such as \( t.a = r.x \text{ AND } t.b = r.y \) cannot select more than one \( r \) row per \( t \) row. The planner formerly considered these \texttt{AND} conditions to be independent and would often drastically misestimate selectivity as a result. Now it compares the \texttt{WHERE} conditions to applicable foreign key constraints and produces better estimates.

E.18.3.1.6. \texttt{VACUUM}

- Avoid re-vacuuming pages containing only frozen tuples (Masahiko Sawada, Robert Haas, Andres Freund)

Formerly, anti-wraparound vacuum had to visit every page of a table, even pages where there was nothing to do. Now, pages containing only already-frozen tuples are identified in the table’s visibility map, and can be skipped by vacuum even when doing transaction wraparound prevention. This should greatly reduce the cost of maintaining large tables containing mostly-unchanging data.

If necessary, vacuum can be forced to process all-frozen pages using the new \texttt{DISABLE_PAGE_SKIPPING} option. Normally this should never be needed, but it might help in recovering from visibility-map corruption.

- Avoid useless heap-truncation attempts during \texttt{VACUUM} (Jeff Janes, Tom Lane)

This change avoids taking an exclusive table lock in some cases where no truncation is possible. The main benefit comes from avoiding unnecessary query cancellations on standby servers.

E.18.3.1.7. General Performance

- Allow old MVCC snapshots to be invalidated after a configurable timeout (Kevin Grittner)

Normally, deleted tuples cannot be physically removed by vacuuming until the last transaction that could “see” them is gone. A transaction that stays open for a long time can thus cause considerable table bloat because space cannot be recycled. This feature allows setting a time-based limit, via the new configuration parameter \texttt{old_snapshot_threshold}, on how long an MVCC snapshot is guaranteed to be valid. After that, dead tuples are candidates for removal. A transaction using an outdated snapshot will get an error if it attempts to read a page that potentially could have contained such data.

- Ignore \texttt{GROUP BY} columns that are functionally dependent on other columns (David Rowley)

If a \texttt{GROUP BY} clause includes all columns of a non-deferred primary key, as well as other columns of the same table, those other columns are redundant and can be dropped from the grouping. This saves computation in many common cases.

- Allow use of an index-only scan on a partial index when the index’s \texttt{WHERE} clause references columns that are not indexed (Tomas Vondra, Kyotaro Horiguchi)

For example, an index defined by \texttt{CREATE INDEX tidx\_partial ON t(b) WHERE a > 0} can now be used for an index-only scan by a query that specifies \texttt{WHERE a > 0} and does not otherwise use \( a \). Previously this was disallowed because \( a \) is not listed as an index column.

- Perform checkpoint writes in sorted order (Fabien Coelho, Andres Freund)

Previously, checkpoints wrote out dirty pages in whatever order they happen to appear in shared buffers, which usually is nearly random. That performs poorly, especially on rotating media. This change causes checkpoint-driven writes to be done in order by file and block number, and to be balanced across tablespaces.
Appendix E. Release Notes

- Where feasible, trigger kernel writeback after a configurable number of writes, to prevent accumulation of dirty data in kernel disk buffers (Fabien Coelho, Andres Freund)

PostgreSQL writes data to the kernel’s disk cache, from where it will be flushed to physical storage in due time. Many operating systems are not smart about managing this and allow large amounts of dirty data to accumulate before deciding to flush it all at once, causing long delays for new I/O requests until the flushing finishes. This change attempts to alleviate this problem by explicitly requesting data flushes after a configurable interval.

On Linux, `sync_file_range()` is used for this purpose, and the feature is on by default on Linux because that function has few downsides. This flushing capability is also available on other platforms if they have `msync()` or `posix_fadvise()`, but those interfaces have some undesirable side-effects so the feature is disabled by default on non-Linux platforms.

The new configuration parameters `backend_flush_after`, `bgwriter_flush_after`, `checkpoint_flush_after`, and `wal_writer_flush_after` control this behavior.

- Improve aggregate-function performance by sharing calculations across multiple aggregates if they have the same arguments and transition functions (David Rowley)

For example, `SELECT AVG(x), VARIANCE(x) FROM tab` can use a single per-row computation for both aggregates.

- Speed up visibility tests for recently-created tuples by checking the current transaction’s snapshot, not `pg_clog`, to decide if the source transaction should be considered committed (Jeff Janes, Tom Lane)

- Allow tuple hint bits to be set sooner than before (Andres Freund)

- Improve performance of short-lived prepared transactions (Stas Kelvich, Simon Riggs, Pavan Deolasee)

Two-phase commit information is now written only to WAL during `PREPARE TRANSACTION`, and will be read back from WAL during `COMMIT PREPARED` if that happens soon thereafter. A separate state file is created only if the pending transaction does not get committed or aborted by the time of the next checkpoint.

- Improve performance of memory context destruction (Jan Wieck)

- Improve performance of resource owners with many tracked objects (Aleksander Alekseev)

- Improve speed of the output functions for `timestamp`, `time`, and `date` data types (David Rowley, Andres Freund)

- Avoid some unnecessary cancellations of hot-standby queries during replay of actions that take `AccessExclusive` locks (Jeff Janes)

- Extend relations multiple blocks at a time when there is contention for the relation’s extension lock (Dilip Kumar)

This improves scalability by decreasing contention.

- Increase the number of clog buffers for better scalability (Amit Kapila, Andres Freund)

- Speed up expression evaluation in PL/pgSQL by keeping `ParamListInfo` entries for simple variables valid at all times (Tom Lane)

- Avoid reducing the `SO_SNDBUF` setting below its default on recent Windows versions (Chen Hua-jun)

- Disable `update_process_title` by default on Windows (Takayuki Tsunakawa)
Appendix E. Release Notes

The overhead of updating the process title is much larger on Windows than most other platforms, and it is also less useful to do it since most Windows users do not have tools that can display process titles.

E.18.3.1.8. Monitoring

- Add `pg_stat_progress_vacuum` system view to provide progress reporting for VACUUM operations (Amit Langote, Robert Haas, Vinayak Pokale, Rahila Syed)
- Add `pg_control_system()`, `pg_control_checkpoint()`, `pg_control_recovery()`, and `pg_control_init()` functions to expose fields of `pg_control` to SQL (Joe Conway, Michael Paquier)
- Add `pg_config` system view (Joe Conway)
  This view exposes the same information available from the `pg_config` command-line utility, namely assorted compile-time configuration information for PostgreSQL.
- Add a `confirmed_flush_lsn` column to the `pg_replication_slots` system view (Marko Tiikkaja)
- Add `pg_stat_wal_receiver` system view to provide information about the state of a hot-standby server’s WAL receiver process (Michael Paquier)
- Add `pg_blocking_pids()` function to reliably identify which sessions block which others (Tom Lane)
  This function returns an array of the process IDs of any sessions that are blocking the session with the given process ID. Historically users have obtained such information using a self-join on the `pg_locks` view. However, it is unreasonably tedious to do it that way with any modicum of correctness, and the addition of parallel queries has made the old approach entirely impractical, since locks might be held or awaited by child worker processes rather than the session’s main process.
- Add function `pg_current_xlog_flush_location()` to expose the current transaction log flush location (Tomas Vondra)
- Add function `pg_notification_queue_usage()` to report how full the NOTIFY queue is (Brendan Jurd)
- Limit the verbosity of memory context statistics dumps (Tom Lane)
  The memory usage dump that is output to the postmaster log during an out-of-memory failure now summarizes statistics when there are a large number of memory contexts, rather than possibly generating a very large report. There is also a “grand total” summary line now.

E.18.3.1.9. Authentication

- Add a BSD authentication method to allow use of the BSD Authentication service for PostgreSQL client authentication (Marisa Emerson)
  BSD Authentication is currently only available on OpenBSD.
- When using PAM authentication, provide the client IP address or host name to PAM modules via the `PAM_RHOST` item (Grzegorz Sampolski)
- Provide detail in the postmaster log for more types of password authentication failure (Tom Lane)
All ordinarily-reachable password authentication failure cases should now provide specific `DETAIL` fields in the log.

- Support RADIUS passwords up to 128 characters long (Marko Tiikkaja)
- Add new SSPI authentication parameters `compat_realm` and `upn_username` to control whether NetBIOS or Kerberos realm names and user names are used during SSPI authentication (Christian Ullrich)

**E.18.3.1.10. Server Configuration**

- Allow sessions to be terminated automatically if they are in idle-in-transaction state for too long (Vik Fearing)
  
  This behavior is controlled by the new configuration parameter `idle_in_transaction_session_timeout`. It can be useful to prevent forgotten transactions from holding locks or preventing vacuum cleanup for too long.
- Raise the maximum allowed value of `checkpoint_timeout` to 24 hours (Simon Riggs)
- Allow `effective_io_concurrency` to be set per-tablespace to support cases where different tablespaces have different I/O characteristics (Julien Rouhaud)
- Add `log_line_prefix` option `&n` to print the current time in Unix epoch form, with milliseconds (Tomas Vondra, Jeff Davis)
- Add `syslog_sequence_numbers` and `syslog_split_messages` configuration parameters to provide more control over the message format when logging to syslog (Peter Eisentraut)
- Merge the `archive` and `hot_standby` values of the `wal_level` configuration parameter into a single new value `replica` (Peter Eisentraut)
  
  Making a distinction between these settings is no longer useful, and merging them is a step towards a planned future simplification of replication setup. The old names are still accepted but are converted to `replica` internally.
- Add configure option `--with-systemd` to enable calling `sd_notify()` at server start and stop (Peter Eisentraut)
  
  This allows the use of systemd service units of type `notify`, which greatly simplifies the management of PostgreSQL under systemd.
- Allow the server’s SSL key file to have group read access if it is owned by `root` (Christoph Berg)
  
  Formerly, we insisted the key file be owned by the user running the PostgreSQL server, but that is inconvenient on some systems (such as Debian) that are configured to manage certificates centrally. Therefore, allow the case where the key file is owned by `root` and has group read access. It is up to the operating system administrator to ensure that the group does not include any untrusted users.

**E.18.3.1.11. Reliability**

- Force backends to exit if the postmaster dies (Rajeev Rastogi, Robert Haas)
  
  Under normal circumstances the postmaster should always outlive its child processes. If for some reason the postmaster dies, force backend sessions to exit with an error. Formerly, existing backends would continue to run until their clients disconnect, but that is unsafe and inefficient. It also prevents a new postmaster from being started until the last old backend has exited. Backends will detect
postmaster death when waiting for client I/O, so the exit will not be instantaneous, but it should happen no later than the end of the current query.

- Check for serializability conflicts before reporting constraint-violation failures (Thomas Munro)

When using serializable transaction isolation, it is desirable that any error due to concurrent transactions should manifest as a serialization failure, thereby cueing the application that a retry might succeed. Unfortunately, this does not reliably happen for duplicate-key failures caused by concurrent insertions. This change ensures that such an error will be reported as a serialization error if the application explicitly checked for the presence of a conflicting key (and did not find it) earlier in the transaction.

- Ensure that invalidation messages are recorded in WAL even when issued by a transaction that has no XID assigned (Andres Freund)

This fixes some corner cases in which transactions on standby servers failed to notice changes, such as new indexes.

- Prevent multiple processes from trying to clean a GIN index’s pending list concurrently (Teodor Sigaev, Jeff Janes)

This had been intentionally allowed, but it causes race conditions that can result in vacuum missing index entries it needs to delete.

**E.18.3.2. Replication and Recovery**

- Allow synchronous replication to support multiple simultaneous synchronous standby servers, not just one (Masahiko Sawada, Beena Emerson, Michael Paquier, Fujii Masao, Kyotaro Horiguchi)

The number of standby servers that must acknowledge a commit before it is considered complete is now configurable as part of the synchronous_standby_names parameter.

- Add new setting `remote_apply` for configuration parameter `synchronous_commit` (Thomas Munro)

In this mode, the master waits for the transaction to be applied on the standby server, not just written to disk. That means that you can count on a transaction started on the standby to see all commits previously acknowledged by the master.

- Add a feature to the replication protocol, and a corresponding option to `pg_create_physical_replication_slot()`, to allow reserving WAL immediately when creating a replication slot (Gurjeet Singh, Michael Paquier)

This allows the creation of a replication slot to guarantee that all the WAL needed for a base backup will be available.

- Add a `--slot` option to `pg_basebackup` (Peter Eisentraut)

This lets `pg_basebackup` use a replication slot defined for WAL streaming. After the base backup completes, selecting the same slot for regular streaming replication allows seamless startup of the new standby server.

- Extend `pg_start_backup()` and `pg_stop_backup()` to support non-exclusive backups (Magnus Hagander)
E.18.3.3. Queries

- Allow functions that return sets of tuples to return simple NULLs (Andrew Gierth, Tom Lane)
  
  In the context of SELECT FROM function(...), a function that returned a set of composite values was previously not allowed to return a plain NULL value as part of the set. Now that is allowed and interpreted as a row of NULLs. This avoids corner-case errors with, for example, unnesting an array of composite values.

- Fully support array subscripts and field selections in the target column list of an INSERT with multiple VALUES rows (Tom Lane)
  
  Previously, such cases failed if the same target column was mentioned more than once, e.g., INSERT INTO tab (x[1], x[2]) VALUES (...).

- When appropriate, postpone evaluation of SELECT output expressions until after an ORDER BY sort (Konstantin Knizhnik)
  
  This change ensures that volatile or expensive functions in the output list are executed in the order suggested by ORDER BY, and that they are not evaluated more times than required when there is a LIMIT clause. Previously, these properties held if the ordering was performed by an index scan or pre-merge-join sort, but not if it was performed by a top-level sort.

- Widen counters recording the number of tuples processed to 64 bits (Andreas Scherbaum)
  
  This change allows command tags, e.g., SELECT, to correctly report tuple counts larger than 4 billion. This also applies to PL/pgSQL's GET DIAGNOSTICS ... ROW_COUNT command.

- Avoid doing encoding conversions by converting through the MULE_INTERNAL encoding (Tom Lane)
  
  Previously, many conversions for Cyrillic and Central European single-byte encodings were done by converting to a related MULE_INTERNAL coding scheme and then to the destination encoding. Aside from being inefficient, this meant that when the conversion encountered an untranslatable character, the error message would confusingly complain about failure to convert to or from MULE_INTERNAL, rather than the user-visible encoding.

- Consider performing joins of foreign tables remotely only when the tables will be accessed under the same role ID (Shigeru Hanada, Ashutosh Bapat, Etsuro Fujita)
  
  Previously, the foreign join pushdown infrastructure left the question of security entirely up to individual foreign data wrappers, but that made it too easy for an FDW to inadvertently create subtle security holes. So, make it the core code’s job to determine which role ID will access each table, and do not attempt join pushdown unless the role is the same for all relevant relations.

E.18.3.4. Utility Commands

- Allow COPY to copy the output of an INSERT/UPDATE/DELETE ... RETURNING query (Marko Tiikkaja)
  
  Previously, an intermediate CTE had to be written to get this result.

- Introduce ALTER object DEPENDS ON EXTENSION (Abhijit Menon-Sen)
  
  This command allows a database object to be marked as depending on an extension, so that it will be dropped automatically if the extension is dropped (without needing CASCADE). However, the object is not part of the extension, and thus will be dumped separately by pg_dump.
• Make `ALTER object SET SCHEMA` do nothing when the object is already in the requested schema, rather than throwing an error as it historically has for most object types (Marti Raudsepp)

• Add options to `ALTER OPERATOR` to allow changing the selectivity functions associated with an existing operator (Yury Zhuravlev)

• Add an `IF NOT EXISTS` option to `ALTER TABLE ADD COLUMN` (Fabrizio de Royes Mello)

• Reduce the lock strength needed by `ALTER TABLE` when setting fillfactor and autovacuum-related relation options (Fabrizio de Royes Mello, Simon Riggs)

• Introduce `CREATE ACCESS METHOD` to allow extensions to create index access methods (Alexander Korotkov, Petr Jelinek)

• Add a `CASCADE` option to `CREATE EXTENSION` to automatically create any extensions the requested one depends on (Petr Jelinek)

• Make `CREATE TABLE ... LIKE` include an `OID` column if any source table has one (Bruce Momjian)

• If a `CHECK` constraint is declared `NOT VALID` in a table creation command, automatically mark it as valid (Amit Langote, Amul Sul)

  This is safe because the table has no existing rows. This matches the longstanding behavior of `FOREIGN KEY` constraints.

• Fix `DROP OPERATOR` to clear `pg_operator.oprcom` and `pg_operator.oprnegate` links to the dropped operator (Roma Sokolov)

  Formerly such links were left as-is, which could pose a problem in the somewhat unlikely event that the dropped operator’s OID was reused for another operator.

• Do not show the same subplan twice in `EXPLAIN` output (Tom Lane)

  In certain cases, typically involving SubPlan nodes in index conditions, `EXPLAIN` would print data for the same subplan twice.

• Disallow creation of indexes on system columns, except for `OID` columns (David Rowley)

  Such indexes were never considered supported, and would very possibly misbehave since the system might change the system-column fields of a tuple without updating indexes. However, previously there were no error checks to prevent them from being created.

### E.18.3.5. Permissions Management

• Use the privilege system to manage access to sensitive functions (Stephen Frost)

  Formerly, many security-sensitive functions contained hard-wired checks that would throw an error if they were called by a non-superuser. This forced the use of superuser roles for some relatively pedestrian tasks. The hard-wired error checks are now gone in favor of making `initdb` revoke the default public `EXECUTE` privilege on these functions. This allows installations to choose to grant usage of such functions to trusted roles that do not need all superuser privileges.

• Create some built-in roles that can be used to grant access to what were previously superuser-only functions (Stephen Frost)

  Currently the only such role is `pg_signal_backend`, but more are expected to be added in future.
E.18.3.6. Data Types

- Improve full-text search to support searching for phrases, that is, lexemes appearing adjacent to each other in a specific order, or with a specified distance between them (Teodor Sigaev, Oleg Bartunov, Dmitry Ivanov)

A phrase-search query can be specified in tsquery input using the new operators \(<\sim\>\) and \(<N>\). The former means that the lexemes before and after it must appear adjacent to each other in that order. The latter means they must be exactly \(N\) lexemes apart.

- Allow omitting one or both boundaries in an array slice specifier, e.g. `array_col[3:]` (Yury Zhuravlev)

Omitted boundaries are taken as the upper or lower limit of the corresponding array subscript. This allows simpler specification for many common use-cases.

- Be more careful about out-of-range dates and timestamps (Vitaly Burovoy)

This change prevents unexpected out-of-range errors for `timestamp with time zone` values very close to the implementation limits. Previously, the “same” value might be accepted or not depending on the `timezone` setting, meaning that a dump and reload could fail on a value that had been accepted when presented. Now the limits are enforced according to the equivalent UTC time, not local time, so as to be independent of `timezone`.

Also, PostgreSQL is now more careful to detect overflow in operations that compute new date or timestamp values, such as `date + integer`.

- For geometric data types, make sure `infinity` and `NaN` component values are treated consistently during input and output (Tom Lane)

Such values will now always print the same as they would in a simple `float8` column, and be accepted the same way on input. Previously the behavior was platform-dependent.

- Upgrade the `ispell` dictionary type to handle modern Hunspell files and support more languages (Artur Zakirov)

- Implement look-behind constraints in regular expressions (Tom Lane)

A look-behind constraint is like a lookahead constraint in that it consumes no text; but it checks for existence (or nonexistence) of a match ending at the current point in the string, rather than one starting at the current point. Similar features exist in many other regular-expression engines.

- In regular expressions, if an apparent three-digit octal escape `\nnn` would exceed 377 (255 decimal), assume it is a two-digit octal escape instead (Tom Lane)

This makes the behavior match current Tcl releases.

- Add transaction ID operators `xid <\sim\> xid` and `xid <\sim\> int4`, for consistency with the corresponding equality operators (Michael Paquier)

E.18.3.7. Functions

- Add `jsonb_insert()` function to insert a new element into a `jsonb` array, or a not-previously-existing key into a `jsonb` object (Dmitry Dolgov)

- Improve the accuracy of the `ln()`, `log()`, `exp()`, and `pow()` functions for type `numeric` (Dean Rasheed)

- Add a `scale(numeric)` function to extract the display scale of a numeric value (Marko Tiikkaja)
Appendix E. Release Notes

- Add trigonometric functions that work in degrees (Dean Rasheed)
  For example, \( \text{sind()} \) measures its argument in degrees, whereas \( \sin() \) measures in radians. These functions go to some lengths to deliver exact results for values where an exact result can be expected, for instance \( \text{sind}(30) = 0.5 \).

- Ensure that trigonometric functions handle infinity and NaN inputs per the POSIX standard (Dean Rasheed)
  The POSIX standard says that these functions should return NaN for NaN input, and should throw an error for out-of-range inputs including infinity. Previously our behavior varied across platforms.

- Make \( \text{to_timestamp(float8)} \) convert float infinity to timestamp infinity (Vitaly Burovoy)
  Formerly it just failed on an infinite input.

- Add new functions for \text{tsvector} data (Stas Kelvich)
  The new functions are \( \text{ts_delete()}, \text{ts_filter()}, \text{unnest()}, \text{tsvector_to_array()}, \text{array_to_tsvector}(), \) and a variant of \text{setweight()} that sets the weight only for specified lexeme(s).

- Allow \( \text{ts_stat()} \) and \( \text{tsvector_update_trigger()} \) to operate on values that are of types binary-compatible with the expected argument type, not just exactly that type; for example allow \text{citext} where \text{text} is expected (Teodor Sigaev)

- Add variadic functions \( \text{num_nulls()} \) and \( \text{num_nonnulls()} \) that count the number of their arguments that are null or non-null (Marko Tiikkaja)
  An example usage is \( \text{CHECK(num_nonnulls(a,b,c) = 1)} \) which asserts that exactly one of \( a,b,c \) is not NULL. These functions can also be used to count the number of null or nonnull elements in an array.

- Add function \( \text{parse_ident()} \) to split a qualified, possibly quoted SQL identifier into its parts (Pavel Stehule)

- In \( \text{to_number()} \), interpret a \( V \) format code as dividing by 10 to the power of the number of digits following \( V \) (Bruce Momjian)
  This makes it operate in an inverse fashion to \( \text{to_char()} \).

- Make the \( \text{to_reg*()} \) functions accept type \text{text} not \text{cstring} (Petr Korobeinikov)
  This avoids the need to write an explicit cast in most cases where the argument is not a simple literal constant.

- Add \( \text{pg_size_bytes()} \) function to convert human-readable size strings to numbers (Pavel Stehule, Vitaly Burovoy, Dean Rasheed)
  This function converts strings like those produced by \( \text{pg_size_pretty()} \) into bytes. An example usage is \( \text{SELECT oid::regclass FROM pg_class WHERE pg_total_relation_size(oid) > pg_size_bytes('10 GB')} \).

- In \( \text{pg_size_pretty()} \), format negative numbers similarly to positive ones (Adrian Vondendriesch)
  Previously, negative numbers were never abbreviated, just printed in bytes.

- Add an optional \( \text{missing_ok} \) argument to the \( \text{current_setting()} \) function (David Christensen)
  This allows avoiding an error for an unrecognized parameter name, instead returning a \text{NULL}.

- Change various catalog-inspection functions to return \text{NULL} for invalid input (Michael Paquier)
Appendix E. Release Notes

pg_get_viewdef() now returns NULL if given an invalid view OID, and several similar functions likewise return NULL for bad input. Previously, such cases usually led to "cache lookup failed" errors, which are not meant to occur in user-facing cases.

- Fix pg_replication_origin_xact_reset() to not have any arguments (Fujii Masao)
  The documentation said that it has no arguments, and the C code did not expect any arguments, but the entry in pg_proc mistakenly specified two arguments.

E.18.3.8. Server-Side Languages

- In PL/pgSQL, detect mismatched CONTINUE and EXIT statements while compiling a function, rather than at execution time (Jim Nasby)
- Extend PL/Python’s error-reporting and message-reporting functions to allow specifying additional message fields besides the primary error message (Pavel Stehule)
- Allow PL/Python functions to call themselves recursively via SPI, and fix the behavior when multiple set-returning PL/Python functions are called within one query (Alexey Grishchenko, Tom Lane)
- Fix session-lifespan memory leaks in PL/Python (Heikki Linnakangas, Haribabu Kommi, Tom Lane)
- Modernize PL/Tcl to use Tcl’s "object" APIs instead of simple strings (Jim Nasby, Karl Lehenbauer)
  This can improve performance substantially in some cases. Note that PL/Tcl now requires Tcl 8.4 or later.
- In PL/Tcl, make database-reported errors return additional information in Tcl’s errorCode global variable (Jim Nasby, Tom Lane)
  This feature follows the Tcl convention for returning auxiliary data about an error.
- Fix PL/Tcl to perform encoding conversion between the database encoding and UTF-8, which is what Tcl expects (Tom Lane)
  Previously, strings were passed through without conversion, leading to misbehavior with non-ASCII characters when the database encoding was not UTF-8.

E.18.3.9. Client Interfaces

- Add a nonlocalized version of the severity field in error and notice messages (Tom Lane)
  This change allows client code to determine severity of an error or notice without having to worry about localized variants of the severity strings.
- Introduce a feature in libpq whereby the CONTEXT field of messages can be suppressed, either always or only for non-error messages (Pavel Stehule)
  The default behavior of PQerrorMessage() is now to print CONTEXT only for errors. The new function PQsetErrorContextVisibility() can be used to adjust this.
- Add support in libpq for regenerating an error message with a different verbosity level (Alex Shulgin)
This is done with the new function PQresultVerboseErrorMessage(). This supports psql’s new \errverbose feature, and may be useful for other clients as well.

- Improve libpq’s PQhost() function to return useful data for default Unix-socket connections (Tom Lane)
  Previously it would return NULL if no explicit host specification had been given; now it returns the default socket directory path.

- Fix ecpg’s lexer to handle line breaks within comments starting on preprocessor directive lines (Michael Meskes)

### E.18.3.10. Client Applications

- Add a --strict-names option to pg_dump and pg_restore (Pavel Stehule)
  This option causes the program to complain if there is no match for a -t or -n option, rather than silently doing nothing.

- In pg_dump, dump locally-made changes of privilege assignments for system objects (Stephen Frost)
  While it has always been possible for a superuser to change the privilege assignments for built-in or extension-created objects, such changes were formerly lost in a dump and reload. Now, pg_dump recognizes and dumps such changes. (This works only when dumping from a 9.6 or later server, however.)

- Allow pg_dump to dump non-extension-owned objects that are within an extension-owned schema (Martín Marqués)
  Previously such objects were ignored because they were mistakenly assumed to belong to the extension owning their schema.

- In pg_dump output, include the table name in object tags for object types that are only uniquely named per-table (for example, triggers) (Peter Eisentraut)

### E.18.3.10.1. psql

- Support multiple -c and -f command-line options (Pavel Stehule, Catalin Iacob)
  The specified operations are carried out in the order in which the options are given, and then psql terminates.

- Add a \crosstabview command that prints the results of a query in a cross-tabulated display (Daniel Vérité)
  In the crosstab display, data values from one query result column are placed in a grid whose column and row headers come from other query result columns.

- Add an \errverbose command that shows the last server error at full verbosity (Alex Shulgin)
  This is useful after getting an unexpected error — you no longer need to adjust the VERBOSITY variable and recreate the failure in order to see error fields that are not shown by default.

- Add \ev and \sv commands for editing and showing view definitions (Petr Korobeinikov)
  These are parallel to the existing \ef and \sf commands for functions.

- Add a \gexec command that executes a query and re-submits the result(s) as new queries (Corey Huinker)
Appendix E. Release Notes

- Allow \pset C string to set the table title, for consistency with \C string (Bruce Momjian)
- In \pset expanded auto mode, do not use expanded format for query results with only one column (Andreas Karlsson, Robert Haas)
- Improve the headers output by the \watch command (Michael Paquier, Tom Lane)
  Include the \pset title string if one has been set, and shorten the prefabricated part of the header to be timestamp (every Ns). Also, the timestamp format now obeys psql’s locale environment.
- Improve tab-completion logic to consider the entire input query, not only the current line (Tom Lane)
  Previously, breaking a command into multiple lines defeated any tab completion rules that needed to see words on earlier lines.
- Numerous minor improvements in tab-completion behavior (Peter Eisentraut, Vik Fearing, Kevin Grittner, Kyotaro Horiguchi, Jeff Janes, Andreas Karlsson, Fujii Masao, Thomas Munro, Masahiko Sawada, Pavel Stehule)
- Add a PROMPT option %p to insert the process ID of the connected backend (Julien Rouhaud)
- Introduce a feature whereby the CONTEXT field of messages can be suppressed, either always or only for non-error messages (Pavel Stehule)
  Printing CONTEXT only for errors is now the default behavior. This can be changed by setting the special variable SHOW_CONTEXT.
- Make \df+ show function access privileges and parallel-safety attributes (Michael Paquier)

E.18.3.10.2. pgbench

- SQL commands in pgbench scripts are now ended by semicolons, not newlines (Kyotaro Horiguchi, Tom Lane)
  This change allows SQL commands in scripts to span multiple lines. Existing custom scripts will need to be modified to add a semicolon at the end of each line that does not have one already. (Doing so does not break the script for use with older versions of pgbench.)
- Support floating-point arithmetic, as well as some built-in functions, in expressions in backslash commands (Fabien Coelho)
- Replace \setrandom with built-in functions (Fabien Coelho)
  The new built-in functions include random(), random_exponential(), and random_gaussian(), which perform the same work as \setrandom, but are easier to use since they can be embedded in larger expressions. Since these additions have made \setrandom obsolete, remove it.
- Allow invocation of multiple copies of the built-in scripts, not only custom scripts (Fabien Coelho)
  This is done with the new -b switch, which works similarly to -f for custom scripts.
- Allow changing the selection probabilities (weights) for scripts (Fabien Coelho)
  When multiple scripts are specified, each pgbench transaction randomly chooses one to execute. Formerly this was always done with uniform probability, but now different selection probabilities can be specified for different scripts.
- Collect statistics for each script in a multi-script run (Fabien Coelho)
  This feature adds an intermediate level of detail to existing global and per-command statistics printouts.
Appendix E. Release Notes

- Add a --progress-timestamp option to report progress with Unix epoch timestamps, instead of time since the run started (Fabien Coelho)
- Allow the number of client connections (-c) to not be an exact multiple of the number of threads (-j) (Fabien Coelho)
- When the -T option is used, stop promptly at the end of the specified time (Fabien Coelho)
  Previously, specifying a low transaction rate could cause pgbench to wait significantly longer than specified.

E.18.3.11. Server Applications

- Improve error reporting during initdb’s post-bootstrap phase (Tom Lane)
  Previously, an error here led to reporting the entire input file as the “failing query”; now just the current query is reported. To get the desired behavior, queries in initdb’s input files must be separated by blank lines.
- Speed up initdb by using just one standalone-backend session for all the post-bootstrap steps (Tom Lane)
- Improve pg_rewind so that it can work when the target timeline changes (Alexander Korotkov)
  This allows, for example, rewinding a promoted standby back to some state of the old master’s timeline.

E.18.3.12. Source Code

- Remove obsolete heap_formtuple/heap_modifytuple/heap_deformtuple functions (Peter Geoghegan)
- Add macros to make AllocSetContextCreate() calls simpler and safer (Tom Lane)
  Writing out the individual sizing parameters for a memory context is now deprecated in favor of using one of the new macros ALLOCSET_DEFAULT_SIZES, ALLOCSET_SMALL_SIZES, or ALLOCSET_START_SMALL_SIZES. Existing code continues to work, however.
- Unconditionally use static inline functions in header files (Andres Freund)
  This may result in warnings and/or wasted code space with very old compilers, but the notational improvement seems worth it.
- Improve TAP testing infrastructure (Michael Paquier, Craig Ringer, Álvaro Herrera, Stephen Frost)
  Notably, it is now possible to test recovery scenarios using this infrastructure.
- Make trace_lwlocks identify individual locks by name (Robert Haas)
- Improve psql’s tab-completion code infrastructure (Thomas Munro, Michael Paquier)
  Tab-completion rules are now considerably easier to write, and more compact.
- Nail the pg_shseclabel system catalog into cache, so that it is available for access during connection authentication (Adam Brightwell)
  The core code does not use this catalog for authentication, but extensions might wish to consult it.
Appendix E. Release Notes

- Restructure index access method API to hide most of it at the C level (Alexander Korotkov, Andrew Gierth)
  This change modernizes the index AM API to look more like the designs we have adopted for foreign data wrappers and tablesample handlers. This simplifies the C code and makes it much more practical to define index access methods in installable extensions. A consequence is that most of the columns of the `pg_am` system catalog have disappeared. New inspection functions have been added to allow SQL queries to determine index AM properties that used to be discoverable from `pg_am`.

- Add `pg_init_privs` system catalog to hold original privileges of initdb-created and extension-created objects (Stephen Frost)
  This infrastructure allows `pg_dump` to dump changes that an installation may have made in privileges attached to system objects. Formerly, such changes would be lost in a dump and reload, but now they are preserved.

- Change the way that extensions allocate custom LWLocks (Amit Kapila, Robert Haas)
  The `RequestAddinLWLocks()` function is removed, and replaced by `RequestNamedLWLockTranche()`. This allows better identification of custom LWLocks, and is less error-prone.

- Improve the isolation tester to allow multiple sessions to wait concurrently, allowing testing of deadlock scenarios (Robert Haas)

- Introduce extensible node types (KaiGai Kohei)
  This change allows FDWs or custom scan providers to store data in a plan tree in a more convenient format than was previously possible.

- Make the planner deal with post-scan/join query steps by generating and comparing Paths, replacing a lot of ad-hoc logic (Tom Lane)
  This change provides only marginal user-visible improvements today, but it enables future work on a lot of upper-planner improvements that were impractical to tackle using the old code structure.

- Support partial aggregation (David Rowley, Simon Riggs)
  This change allows the computation of an aggregate function to be split into separate parts, for example so that parallel worker processes can cooperate on computing an aggregate. In future it might allow aggregation across local and remote data to occur partially on the remote end.

- Add a generic command progress reporting facility (Vinayak Pokale, Rahila Syed, Amit Langote, Robert Haas)

- Separate out psql’s flex lexer to make it usable by other client programs (Tom Lane, Kyotaro Horiguchi)
  This eliminates code duplication for programs that need to be able to parse SQL commands well enough to identify command boundaries. Doing that in full generality is more painful than one could wish, and up to now only psql has really gotten it right among our supported client programs.
  A new source-code subdirectory `src/fe_utils/` has been created to hold this and other code that is shared across our client programs. Formerly such sharing was accomplished by symbolic linking or copying source files at build time, which was ugly and required duplicate compilation.

- Introduce `WaitEventSet` API to allow efficient waiting for event sets that usually do not change from one wait to the next (Andres Freund, Amit Kapila)

- Add a generic interface for writing WAL records (Alexander Korotkov, Petr Jelínek, Markus Nullmeier)
Appendix E. Release Notes

This change allows extensions to write WAL records for changes to pages using a standard layout. The problem of needing to replay WAL without access to the extension is solved by having generic replay code. This allows extensions to implement, for example, index access methods and have WAL support for them.

- **Support generic WAL messages for logical decoding (Petr Jelínek, Andres Freund)**
  This feature allows extensions to insert data into the WAL stream that can be read by logical-decoding plugins, but is not connected to physical data restoration.

- **Allow SP-GiST operator classes to store an arbitrary “traversal value” while descending the index (Alexander Lebedev, Teodor Sigaev)**
  This is somewhat like the “reconstructed value”, but it could be any arbitrary chunk of data, not necessarily of the same data type as the indexed column.

- **Introduce a **LOG_SERVER_ONLY** message level for ereport() (David Steele)**
  This level acts like **LOG** except that the message is never sent to the client. It is meant for use in auditing and similar applications.

- **Provide a Makefile target to build all generated headers (Michael Paquier, Tom Lane)**
  submake-generated-headers can now be invoked to ensure that generated backend header files are up-to-date. This is useful in subdirectories that might be built “standalone”.

- **Support OpenSSL 1.1.0 (Andreas Karlsson, Heikki Linnakangas)**

E.18.3.13. Additional Modules

- **Add configuration parameter auto_explain.sample_rate to allow contrib/auto_explain to capture just a configurable fraction of all queries (Craig Ringer, Julien Rouhaud)**
  This allows reduction of overhead for heavy query traffic, while still getting useful information on average.

- **Add contrib/bloom module that implements an index access method based on Bloom filtering (Teodor Sigaev, Alexander Korotkov)**
  This is primarily a proof-of-concept for non-core index access methods, but it could be useful in its own right for queries that search many columns.

- **In contrib/cube, introduce distance operators for cubes, and support kNN-style searches in GiST indexes on cube columns (Stas Kelvich)**

- **Make contrib/hstore’s hstore_to_jsonb_loose() and hstore_to_json_loose() functions agree on what is a number (Tom Lane)**
  Previously, hstore_to_jsonb_loose() would convert numeric-looking strings to JSON numbers, rather than strings, even if they did not exactly match the JSON syntax specification for numbers. This was inconsistent with hstore_to_json_loose(), so tighten the test to match the JSON syntax.

- **Add selectivity estimation functions for contrib/intarray operators to improve plans for queries using those operators (Yury Zhuravlev, Alexander Korotkov)**

- **Make contrib/pageinspect’s heap_page_items() function show the raw data in each tuple, and add new functions tuple_data_split() and heap_page_item_attrs() for inspection of individual tuple fields (Nikolay Shaplov)**
• Add an optional S2K iteration count parameter to contrib/pgcrypto’s pgp_sym_encrypt() function (Jeff Janes)

• Add support for “word similarity” to contrib/pg_trgm (Alexander Korotkov, Artur Zakirov)
  These functions and operators measure the similarity between one string and the most similar single word of another string.

• Add configuration parameter pg_trgm.similarity_threshold for contrib/pg_trgm's similarity threshold (Artur Zakirov)
  This threshold has always been configurable, but formerly it was controlled by special-purpose functions set_limit() and show_limit(). Those are now deprecated.

• Improve contrib/pg_trgm's GIN operator class to speed up index searches in which both common and rare keys appear (Jeff Janes)

• Improve performance of similarity searches in contrib/pg_trgm GIN indexes (Christophe Fornaroli)

• Add contrib/pg_visibility module to allow examining table visibility maps (Robert Haas)

• Add ssl_extension_info() function to contrib/sslinfo, to print information about SSL extensions present in the X509 certificate used for the current connection (Dmitry Voronin)

### E.18.3.13.1. postgres_fdw

• Allow extension-provided operators and functions to be sent for remote execution, if the extension is whitelisted in the foreign server’s options (Paul Ramsey)
  Users can enable this feature when the extension is known to exist in a compatible version in the remote database. It allows more efficient execution of queries involving extension operators.

• Consider performing sorts on the remote server (Ashutosh Bapat)

• Consider performing joins on the remote server (Shigeru Hanada, Ashutosh Bapat)

• When feasible, perform UPDATE or DELETE entirely on the remote server (Etsuro Fujita)
  Formerly, remote updates involved sending a SELECT FOR UPDATE command and then updating or deleting the selected rows one-by-one. While that is still necessary if the operation requires any local processing, it can now be done remotely if all elements of the query are safe to send to the remote server.

• Allow the fetch size to be set as a server or table option (Corey Huinker)
  Formerly, postgres_fdw always fetched 100 rows at a time from remote queries; now that behavior is configurable.

• Use a single foreign-server connection for local user IDs that all map to the same remote user (Ashutosh Bapat)

• Transmit query cancellation requests to the remote server (Michael Paquier, Etsuro Fujita)
  Previously, a local query cancellation request did not cause an already-sent remote query to terminate early.
E.19. Prior Releases

Release notes for prior release branches can be found at https://www.postgresql.org/docs/release/
Appendix F. Additional Supplied Modules

This appendix and the next one contain information regarding the modules that can be found in the contrib directory of the PostgreSQL distribution. These include porting tools, analysis utilities, and plug-in features that are not part of the core PostgreSQL system, mainly because they address a limited audience or are too experimental to be part of the main source tree. This does not preclude their usefulness.

This appendix covers extensions and other server plug-in modules found in contrib. Appendix G covers utility programs.

When building from the source distribution, these components are not built automatically, unless you build the "world" target (see step 2). You can build and install all of them by running:

```
make
make install
```

in the contrib directory of a configured source tree; or to build and install just one selected module, do the same in that module’s subdirectory. Many of the modules have regression tests, which can be executed by running:

```
make check
```

before installation or

```
make installcheck
```

once you have a PostgreSQL server running.

If you are using a pre-packaged version of PostgreSQL, these modules are typically made available as a separate subpackage, such as postgresql-contrib.

Many modules supply new user-defined functions, operators, or types. To make use of one of these modules, after you have installed the code you need to register the new SQL objects in the database system. In PostgreSQL 9.1 and later, this is done by executing a CREATE EXTENSION command. In a fresh database, you can simply do

```
CREATE EXTENSION module_name;
```

This command must be run by a database superuser. This registers the new SQL objects in the current database only, so you need to run this command in each database that you want the module’s facilities to be available in. Alternatively, run it in database template1 so that the extension will be copied into subsequently-created databases by default.

Many modules allow you to install their objects in a schema of your choice. To do that, add `SCHEMA schema_name` to the CREATE EXTENSION command. By default, the objects will be placed in your current creation target schema, which in turn defaults to public.

If your database was brought forward by dump and reload from a pre-9.1 version of PostgreSQL, and you had been using the pre-9.1 version of the module in it, you should instead do

```
CREATE EXTENSION module_name FROM unpackaged;
```
This will update the pre-9.1 objects of the module into a proper extension object. Future updates to the module will be managed by ALTER EXTENSION. For more information about extension updates, see Section 36.15.

Note, however, that some of these modules are not “extensions” in this sense, but are loaded into the server in some other way, for instance by way of shared_preload_libraries. See the documentation of each module for details.

F.1. adminpack

adminpack provides a number of support functions which pgAdmin and other administration and management tools can use to provide additional functionality, such as remote management of server log files. Use of all these functions is restricted to superusers.

The functions shown in Table F-1 provide write access to files on the machine hosting the server. (See also the functions in Table 9-86, which provide read-only access.) Only files within the database cluster directory can be accessed, but either a relative or absolute path is allowable.

### Table F-1. adminpack Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_catalog.pg_file_write(filename text, data text, append boolean)</code></td>
<td>bigint</td>
<td>Write, or append to, a text file</td>
</tr>
<tr>
<td><code>pg_catalog.pg_file_rename(oldname text, newname text [, archivename text])</code></td>
<td>boolean</td>
<td>Rename a file</td>
</tr>
<tr>
<td><code>pg_catalog.pg_file_unlink(filename text)</code></td>
<td>boolean</td>
<td>Remove a file</td>
</tr>
<tr>
<td><code>pg_catalog.pg_logdir_ls()</code></td>
<td>setof record</td>
<td>List the log files in the log_directory directory</td>
</tr>
</tbody>
</table>

Pg_file_write writes the specified data into the file named by filename. If append is false, the file must not already exist. If append is true, the file can already exist, and will be appended to if so. Returns the number of bytes written.

Pg_file_rename renames a file. If archivename is omitted or NULL, it simply renames oldname to newname (which must not already exist). If archivename is provided, it first renames newname to archivename (which must not already exist), and then renames oldname to newname. In event of failure of the second rename step, it will try to rename archivename back to newname before reporting the error. Returns true on success, false if the source file(s) are not present or not writable; other cases throw errors.

Pg_file_unlink removes the specified file. Returns true on success, false if the specified file is not present or the unlink() call fails; other cases throw errors.

Pg_logdir_ls returns the start timestamps and path names of all the log files in the log_directory directory. The log_filename parameter must have its default setting (postgresql-%Y-%m-%d_%H%M%S.log) to use this function.

The functions shown in Table F-2 are deprecated and should not be used in new applications; instead use those shown in Table 9-77 and Table 9-86. These functions are provided in adminpack only for compatibility with old versions of pgAdmin.
Table F-2. Deprecated adminpack Functions

<table>
<thead>
<tr>
<th>Name</th>
<th>Return Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>pg_catalog.pg_file_read(filename text, offset bigint, nbytes bigint)</code></td>
<td>text</td>
<td>Alternate name for <code>pg_read_file()</code></td>
</tr>
<tr>
<td><code>pg_catalog.pg_file_length(filename text)</code></td>
<td>bigint</td>
<td>Same as <code>size</code> column returned by <code>pg_stat_file()</code></td>
</tr>
<tr>
<td><code>pg_catalog.pg_logfile_rotate()</code></td>
<td>integer</td>
<td>Alternate name for <code>pg_rotate_logfile()</code>, but note that it returns integer 0 or 1 rather than boolean</td>
</tr>
</tbody>
</table>

F.2. auth_delay

`auth_delay` causes the server to pause briefly before reporting authentication failure, to make brute-force attacks on database passwords more difficult. Note that it does nothing to prevent denial-of-service attacks, and may even exacerbate them, since processes that are waiting before reporting authentication failure will still consume connection slots.

In order to function, this module must be loaded via shared_preload_libraries in `postgresql.conf`.

F.2.1. Configuration Parameters

`auth_delay.milliseconds (int)`

The number of milliseconds to wait before reporting an authentication failure. The default is 0.

These parameters must be set in `postgresql.conf`. Typical usage might be:

```
# postgresql.conf
shared_preload_libraries = 'auth_delay'
auth_delay.milliseconds = '500'
```

F.2.2. Author

KaiGai Kohei <kaigai@ak.jp.nec.com>

F.3. auto_explain

The `auto_explain` module provides a means for logging execution plans of slow statements automatically, without having to run `EXPLAIN` by hand. This is especially helpful for tracking down un-optimized queries in large applications.

The module provides no SQL-accessible functions. To use it, simply load it into the server. You can load it into an individual session:
LOAD 'auto_explain';

(You must be superuser to do that.) More typical usage is to preload it into some or all sessions by including auto_explain in session_preload_libraries or shared_preload_libraries in postgresql.conf. Then you can track unexpectedly slow queries no matter when they happen. Of course there is a price in overhead for that.

**F.3.1. Configuration Parameters**

There are several configuration parameters that control the behavior of auto_explain. Note that the default behavior is to do nothing, so you must set at least auto_explain.log_min_duration if you want any results.

`auto_explain.log_min_duration` (integer)

`auto_explain.log_min_duration` is the minimum statement execution time, in milliseconds, that will cause the statement’s plan to be logged. Setting this to zero logs all plans. Minus-one (the default) disables logging of plans. For example, if you set it to 250ms then all statements that run 250ms or longer will be logged. Only superusers can change this setting.

`auto_explain.log_analyze` (boolean)

`auto_explain.log_analyze` causes EXPLAIN ANALYZE output, rather than just EXPLAIN output, to be printed when an execution plan is logged. This parameter is off by default. Only superusers can change this setting.

**Note:** When this parameter is on, per-plan-node timing occurs for all statements executed, whether or not they run long enough to actually get logged. This can have an extremely negative impact on performance. Turning off `auto_explain.log_timing` ameliorates the performance cost, at the price of obtaining less information.

`auto_explain.log_buffers` (boolean)

`auto_explain.log_buffers` controls whether buffer usage statistics are printed when an execution plan is logged; it’s equivalent to the BUFFERS option of EXPLAIN. This parameter has no effect unless `auto_explain.log_analyze` is enabled. This parameter is off by default. Only superusers can change this setting.

`auto_explain.log_timing` (boolean)

`auto_explain.log_timing` controls whether per-node timing information is printed when an execution plan is logged; it’s equivalent to the TIMING option of EXPLAIN. The overhead of repeatedly reading the system clock can slow down queries significantly on some systems, so it may be useful to set this parameter to off when only actual row counts, and not exact times, are needed. This parameter has no effect unless `auto_explain.log_analyze` is enabled. This parameter is on by default. Only superusers can change this setting.

`auto_explain.log_triggers` (boolean)

`auto_explain.log_triggers` causes trigger execution statistics to be included when an execution plan is logged. This parameter has no effect unless `auto_explain.log_analyze` is enabled. This parameter is off by default. Only superusers can change this setting.
auto_explain.log_verbose (boolean)

auto_explain.log_verbose controls whether verbose details are printed when an execution plan is logged; it’s equivalent to the VERBOSE option of EXPLAIN. This parameter is off by default. Only superusers can change this setting.

auto_explain.log_format (enum)

auto_explain.log_format selects the EXPLAIN output format to be used. The allowed values are text, xml, json, and yaml. The default is text. Only superusers can change this setting.

auto_explain.log_nested_statements (boolean)

auto_explain.log_nested_statements causes nested statements (statements executed inside a function) to be considered for logging. When it is off, only top-level query plans are logged. This parameter is off by default. Only superusers can change this setting.

auto_explain.sample_rate (real)

auto_explain.sample_rate causes auto_explain to only explain a fraction of the statements in each session. The default is 1, meaning explain all the queries. In case of nested statements, either all will be explained or none. Only superusers can change this setting.

In ordinary usage, these parameters are set in postgresql.conf, although superusers can alter them on-the-fly within their own sessions. Typical usage might be:

```sql
# postgresql.conf
session_preload_libraries = 'auto_explain'

auto_explain.log_min_duration = '3s'
```

### F.3.2. Example

postgres=# LOAD 'auto_explain';
postgres=# SET auto_explain.log_min_duration = 0;
postgres=# SET auto_explain.log_analyze = true;
postgres=# SELECT count(*)
  FROM pg_class, pg_index
  WHERE oid = indrelid AND indisunique;

This might produce log output such as:

```
LOG: duration: 3.651 ms plan:
Query Text: SELECT count(*)
  FROM pg_class, pg_index
  WHERE oid = indrelid AND indisunique;
Aggregate  (cost=16.79..16.80 rows=1 width=0) (actual time=3.626..3.627 rows=1 loops=1)
  ->  Hash Join  (cost=4.17..16.55 rows=92 width=0) (actual time=3.349..3.594 rows=92)
    Hash Cond: (pg_class.oid = pg_index.indrelid)
  ->  Seq Scan on pg_class  (cost=0.00..9.55 rows=255 width=4) (actual time=0.016..0.140 rows=255 loops=1)
  ->  Hash  (cost=3.02..3.02 rows=92 width=4) (actual time=3.238..3.238 rows=92)
    Buckets: 1024 Batches: 1 Memory Usage: 4kB
  ->  Seq Scan on pg_index  (cost=0.00..3.02 rows=92 width=4) (actual time=0.008..3.187 rows=92 loops=1)
Filter: indisunique
```
F.3.3. Author

Takahiro Itagaki <itagaki.takahiro@oss.ntt.co.jp>

F.4. bloom

bloom provides an index access method based on Bloom filters¹.

A Bloom filter is a space-efficient data structure that is used to test whether an element is a member of a set. In the case of an index access method, it allows fast exclusion of non-matching tuples via signatures whose size is determined at index creation.

A signature is a lossy representation of the indexed attribute(s), and as such is prone to reporting false positives; that is, it may be reported that an element is in the set, when it is not. So index search results must always be rechecked using the actual attribute values from the heap entry. Larger signatures reduce the odds of a false positive and thus reduce the number of useless heap visits, but of course also make the index larger and hence slower to scan.

This type of index is most useful when a table has many attributes and queries test arbitrary combinations of them. A traditional btree index is faster than a bloom index, but it can require many btree indexes to support all possible queries where one needs only a single bloom index. Note however that bloom indexes only support equality queries, whereas btree indexes can also perform inequality and range searches.

F.4.1. Parameters

A bloom index accepts the following parameters in its WITH clause:

length

Length of each signature (index entry) in bits. It is rounded up to the nearest multiple of 16. The default is 80 bits and the maximum is 4096.

coll -- col32

Number of bits generated for each index column. Each parameter's name refers to the number of the index column that it controls. The default is 2 bits and the maximum is 4095. Parameters for index columns not actually used are ignored.

F.4.2. Examples

This is an example of creating a bloom index:

```sql
CREATE INDEX bloomidx ON tbloom USING bloom (i1,i2,i3)
    WITH (length=80, coll=2, col2=2, col3=4);
```

The index is created with a signature length of 80 bits, with attributes i1 and i2 mapped to 2 bits, and attribute i3 mapped to 4 bits. We could have omitted the length, coll, and col2 specifications since those have the default values.

Here is a more complete example of bloom index definition and usage, as well as a comparison with equivalent btree indexes. The bloom index is considerably smaller than the btree index, and can perform better.

```sql
=# CREATE TABLE tbloom AS
SELECT
  (random() * 1000000)::int as i1,
  (random() * 1000000)::int as i2,
  (random() * 1000000)::int as i3,
  (random() * 1000000)::int as i4,
  (random() * 1000000)::int as i5,
  (random() * 1000000)::int as i6
FROM
  generate_series(1,10000000);
SELECT 10000000
=# CREATE INDEX bloomidx ON tbloom USING bloom (i1, i2, i3, i4, i5, i6);
CREATE INDEX
=# SELECT pg_size_pretty(pg_relation_size('bloomidx'));
pg_size_pretty
----------------
153 MB
(1 row)
=# CREATE index btreeidx ON tbloom (i1, i2, i3, i4, i5, i6);
CREATE INDEX
=# SELECT pg_size_pretty(pg_relation_size('btreeidx'));
pg_size_pretty
----------------
387 MB
(1 row)

A sequential scan over this large table takes a long time:

```sql
=# EXPLAIN ANALYZE SELECT * FROM tbloom WHERE i2 = 898732 AND i5 = 123451;
QUERY PLAN
------------------------------------------------------------------------------------------------------------
Seq Scan on tbloom (cost=0.00..213694.08 rows=1 width=24) (actual time=1445.438..1445.438 rows=0 loops=1)
  Filter: ((i2 = 898732) AND (i5 = 123451))
  Rows Removed by Filter: 10000000
Planning time: 0.177 ms
Execution time: 1445.473 ms
(5 rows)
```

So the planner will usually select an index scan if possible. With a btree index, we get results like this:

```sql
=# EXPLAIN ANALYZE SELECT * FROM tbloom WHERE i2 = 898732 AND i5 = 123451;
QUERY PLAN
--------------------------------------------------------------------------------------------------------------------------------
Index Only Scan using btreeidx on tbloom (cost=0.56..298311.96 rows=1 width=24) (actual time=445.709..445.709 rows=0 loops=1)
  Index Cond: ((i2 = 898732) AND (i5 = 123451))
  Heap Fetches: 0
Planning time: 0.193 ms
Execution time: 445.770 ms
(5 rows)
```
Bloom is better than btree in handling this type of search:

```sql
=# EXPLAIN ANALYZE SELECT * FROM tbloom WHERE i2 = 898732 AND i5 = 123451;
```

```
QUERY PLAN
---------------------------------------------------------------------------------------------------------------------------
  Bitmap Heap Scan on tbloom (cost=178435.39..178439.41 rows=1 width=24) (actual time=76.698..76.698 rows=0 loops=1)
    Recheck Cond: ((i2 = 898732) AND (i5 = 123451))
    Rows Removed by Index Recheck: 2439
    Heap Blocks: exact=2408
  -> Bitmap Index Scan on bloomidx (cost=0.00..178435.39 rows=1 width=0) (actual time=72.455..72.455 rows=2439 loops=1)
      Index Cond: ((i2 = 898732) AND (i5 = 123451))
Planning time: 0.475 ms
Execution time: 76.778 ms
(8 rows)
```

Note the relatively large number of false positives: 2439 rows were selected to be visited in the heap, but none actually matched the query. We could reduce that by specifying a larger signature length. In this example, creating the index with `length=200` reduced the number of false positives to 55; but it doubled the index size (to 306 MB) and ended up being slower for this query (125 ms overall).

Now, the main problem with the btree search is that btree is inefficient when the search conditions do not constrain the leading index column(s). A better strategy for btree is to create a separate index on each column. Then the planner will choose something like this:

```sql
=# EXPLAIN ANALYZE SELECT * FROM tbloom WHERE i2 = 898732 AND i5 = 123451;
```

```
QUERY PLAN
------------------------------------------------------------------------------------------------------------------------------
  Bitmap Heap Scan on tbloom (cost=9.29..13.30 rows=1 width=24) (actual time=0.148..0.148 rows=0 loops=1)
    Recheck Cond: ((i5 = 123451) AND (i2 = 898732))
    -> BitmapAnd (cost=9.29..9.29 rows=1 width=0) (actual time=0.145..0.145 rows=0 loops=1)
       -> Bitmap Index Scan on tbloom_i5_idx (cost=0.00..4.52 rows=11 width=0) (actual time=0.089..0.089 rows=10 loops=1)
          Index Cond: (i5 = 123451)
       -> Bitmap Index Scan on tbloom_i2_idx (cost=0.00..4.52 rows=11 width=0) (actual time=0.048..0.048 rows=8 loops=1)
          Index Cond: (i2 = 898732)
Planning time: 2.049 ms
Execution time: 0.280 ms
(9 rows)
```

Although this query runs much faster than with either of the single indexes, we pay a large penalty in index size. Each of the single-column btree indexes occupies 214 MB, so the total space needed is over 1.2GB, more than 8 times the space used by the bloom index.

### F.4.3. Operator Class Interface

An operator class for bloom indexes requires only a hash function for the indexed data type and an equality operator for searching. This example shows the operator class definition for the `text` data type:

```sql
CREATE OPERATOR CLASS text_ops
DEFAULT FOR TYPE text USING bloom AS
  OPERATOR 1 =(text, text),
  FUNCTION 1 hashtext(text);
```
Appendix F. Additional Supplied Modules

F.4.4. Limitations

- Only operator classes for `int4` and `text` are included with the module.
- Only the `-` operator is supported for search. But it is possible to add support for arrays with union and intersection operations in the future.
- `bloom` access method doesn’t support `UNIQUE` indexes.
- `bloom` access method doesn’t support searching for `NULL` values.

F.4.5. Authors

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Oleg Bartunov <obartunov@postgrespro.ru>, Postgres Professional, Moscow, Russia

F.5. btrees_gin

`btrees_gin` provides sample GIN operator classes that implement B-tree equivalent behavior for the data types `int2`, `int4`, `int8`, `float4`, `float8`, `timestamp` with time zone, `timestamp` without time zone, `time` with time zone, `time` without time zone, `date`, `interval`, `oid`, `money`, "char", `varchar`, `text`, `bytea`, `bit`, `varbit`, `macaddr`, `inet`, and `cidr`.

In general, these operator classes will not outperform the equivalent standard B-tree index methods, and they lack one major feature of the standard B-tree code: the ability to enforce uniqueness. However, they are useful for GIN testing and as a base for developing other GIN operator classes. Also, for queries that test both a GIN-indexable column and a B-tree-indexable column, it might be more efficient to create a multicolumn GIN index that uses one of these operator classes than to create two separate indexes that would have to be combined via bitmap ANDing.

F.5.1. Example Usage

```
CREATE TABLE test (a int4);
-- create index
CREATE INDEX testidx ON test USING GIN (a);
-- query
SELECT * FROM test WHERE a < 10;
```

F.5.2. Authors

Teodor Sigaev (<teodor@stack.net>) and Oleg Bartunov (<oleg@sai.msu.su>). See http://www.sai.msu.su/~megera/oddmuse/index.cgi/Gin for additional information.
Appendix F. Additional Supplied Modules

F.6. btree_gist

**btree_gist** provides GiST index operator classes that implement B-tree equivalent behavior for the data types `int2`, `int4`, `int8`, `float4`, `float8`, `numeric`, `timestamp with time zone`, `time with time zone`, `time without time zone`, `date`, `interval`, `oid`, `money`, `char`, `varchar`, `text`, `bytea`, `bit`, `varbit`, `macaddr`, `inet`, and `cidr`.

In general, these operator classes will not outperform the equivalent standard B-tree index methods, and they lack one major feature of the standard B-tree code: the ability to enforce uniqueness. However, they provide some other features that are not available with a B-tree index, as described below. Also, these operator classes are useful when a multicolored GiST index is needed, wherein some of the columns are of data types that are only indexable with GiST but other columns are just simple data types. Lastly, these operator classes are useful for GiST testing and as a base for developing other GiST operator classes.

In addition to the typical B-tree search operators, **btree_gist** also provides index support for `<>` ("not equals"). This may be useful in combination with an exclusion constraint, as described below.

Also, for data types for which there is a natural distance metric, **btree_gist** defines a distance operator `<=` and provides GiST index support for nearest-neighbor searches using this operator. Distance operators are provided for `int2`, `int4`, `int8`, `float4`, `float8`, `timestamp with time zone`, `time without time zone`, `date`, `interval`, `oid`, and `money`.

### F.6.1. Example Usage

Simple example using **btree_gist** instead of **btree**:

```sql
CREATE TABLE test (a int4);
-- create index
CREATE INDEX testidx ON test USING GIST (a);
-- query
SELECT * FROM test WHERE a < 10;
-- nearest-neighbor search: find the ten entries closest to "42"
SELECT *, a <= 42 AS dist FROM test ORDER BY a <= 42 LIMIT 10;

Use an exclusion constraint to enforce the rule that a cage at a zoo can contain only one kind of animal:

=> CREATE TABLE zoo (cage INTEGER, animal TEXT,
  EXCLUDE USING GIST (cage WITH =, animal WITH <>))

=> INSERT INTO zoo VALUES(123, 'zebra');
INSERT 0 1
=> INSERT INTO zoo VALUES(123, 'zebra');
INSERT 0 1
=> INSERT INTO zoo VALUES(123, 'lion');
ERROR: conflicting key value violates exclusion constraint "zoo_cage_animal_excl"
DETAIL: Key (cage, animal)=(123, lion) conflicts with existing key (cage, animal)=(123,)
=> INSERT INTO zoo VALUES(124, 'lion');
INSERT 0 1
```

2356
F.6.2. Authors

Teodor Sigaev (<teodor@stack.net>), Oleg Bartunov (<oleg@sai.msu.su>), and Janko Richter (<jankorichter@yahoo.de>). See http://www.sai.msu.su/~megera/postgres/gist/ for additional information.

F.7. chkpass

This module implements a data type `chkpass` that is designed for storing encrypted passwords. Each password is automatically converted to encrypted form upon entry, and is always stored encrypted. To compare, simply compare against a clear text password and the comparison function will encrypt it before comparing.

There are provisions in the code to report an error if the password is determined to be easily crackable. However, this is currently just a stub that does nothing.

If you precede an input string with a colon, it is assumed to be an already-encrypted password, and is stored without further encryption. This allows entry of previously-encrypted passwords.

On output, a colon is prepended. This makes it possible to dump and reload passwords without re-encrypting them. If you want the encrypted password without the colon then use the `raw()` function. This allows you to use the type with things like Apache’s Auth_PostgreSQL module.

The encryption uses the standard Unix function `crypt()`, and so it suffers from all the usual limitations of that function; notably that only the first eight characters of a password are considered.

Note that the `chkpass` data type is not indexable.

Sample usage:

```sql
test=# create table test (p chkpass);
CREATE TABLE

test=# insert into test values ('hello');
INSERT 0 1

test=# select * from test;
p
----------------
:dVgpXdxOrE3ko
(1 row)

test=# select raw(p) from test;
raw
-------------
dVgkpXdxOrE3ko
(1 row)

test=# select p = 'hello' from test;
?column?
----------
t
(1 row)

test=# select p = 'goodbye' from test;
?column?
----------
f
```
The citext module provides a case-insensitive character string type, citext. Essentially, it internally calls lower when comparing values. Otherwise, it behaves almost exactly like text.

F.8.1. Rationale

The standard approach to doing case-insensitive matches in PostgreSQL has been to use the lower function when comparing values, for example

```
SELECT * FROM tab WHERE lower(col) = LOWER(?);
```

This works reasonably well, but has a number of drawbacks:

• It makes your SQL statements verbose, and you always have to remember to use lower on both the column and the query value.
• It won’t use an index, unless you create a functional index using lower.
• If you declare a column as UNIQUE or PRIMARY KEY, the implicitly generated index is case-sensitive. So it’s useless for case-insensitive searches, and it won’t enforce uniqueness case-insensitively.

The citext data type allows you to eliminate calls to lower in SQL queries, and allows a primary key to be case-insensitive. citext is locale-aware, just like text, which means that the matching of upper case and lower case characters is dependent on the rules of the database’s LC_CTYPE setting. Again, this behavior is identical to the use of lower in queries. But because it’s done transparently by the data type, you don’t have to remember to do anything special in your queries.

F.8.2. How to Use It

Here’s a simple example of usage:

```sql
CREATE TABLE users (  
nick CITEXT PRIMARY KEY,  
pass TEXT NOT NULL  
);  
INSERT INTO users VALUES ( 'larry', md5(random()::text) );  
INSERT INTO users VALUES ( 'Tom', md5(random()::text) );  
INSERT INTO users VALUES ( 'Damian', md5(random()::text) );  
INSERT INTO users VALUES ( 'NEAL', md5(random()::text) );  
INSERT INTO users VALUES ( 'Bjørn', md5(random()::text) );
```
SELECT * FROM users WHERE nick = 'Larry';

The SELECT statement will return one tuple, even though the nick column was set to larry and the query was for Larry.

F.8.3. String Comparison Behavior

citext performs comparisons by converting each string to lower case (as though lower were called) and then comparing the results normally. Thus, for example, two strings are considered equal if lower would produce identical results for them.

In order to emulate a case-insensitive collation as closely as possible, there are citext-specific versions of a number of string-processing operators and functions. So, for example, the regular expression operators ~ and ~* exhibit the same behavior when applied to citext: they both match case-insensitively. The same is true for !~ and !~*, as well as for the LIKE operators ~* and ~*~, and !~* and !~*~. If you’d like to match case-sensitively, you can cast the operator’s arguments to text.

Similarly, all of the following functions perform matching case-insensitively if their arguments are citext:

• regexp_matches()
• regexp_replace()
• regexp_split_to_array()
• regexp_split_to_table()
• replace()
• split_part()
• strpos()
• translate()

For the regexp functions, if you want to match case-sensitively, you can specify the “c” flag to force a case-sensitive match. Otherwise, you must cast to text before using one of these functions if you want case-sensitive behavior.

F.8.4. Limitations

• citext’s case-folding behavior depends on the LC_CTYPE setting of your database. How it compares values is therefore determined when the database is created. It is not truly case-insensitive in the terms defined by the Unicode standard. Effectively, what this means is that, as long as you’re happy with your collation, you should be happy with citext’s comparisons. But if you have data in different languages stored in your database, users of one language may find their query results are not as expected if the collation is for another language.

• As of PostgreSQL 9.1, you can attach a COLLATE specification to citext columns or data values. Currently, citext operators will honor a non-default COLLATE specification while comparing case-folded strings, but the initial folding to lower case is always done according to the database’s LC_CTYPE setting (that is, as though COLLATE "default" were given). This may be changed in a future release so that both steps follow the input COLLATE specification.
• *citext* is not as efficient as *text* because the operator functions and the B-tree comparison functions must make copies of the data and convert it to lower case for comparisons. It is, however, slightly more efficient than using *lower* to get case-insensitive matching.

• *citext* doesn’t help much if you need data to compare case-sensitively in some contexts and case-insensitively in other contexts. The standard answer is to use the *text* type and manually use the *lower* function when you need to compare case-insensitively; this works all right if case-insensitive comparison is needed only infrequently. If you need case-insensitive behavior most of the time and case-sensitive infrequently, consider storing the data as *citext* and explicitly casting the column to *text* when you want case-sensitive comparison. In either situation, you will need two indexes if you want both types of searches to be fast.

• The schema containing the *citext* operators must be in the current *search_path* (typically *public*); if it is not, the normal case-sensitive *text* operators will be invoked instead.

### F.8.5. Author

David E. Wheeler <david@kineticode.com>

Inspired by the original *citext* module by Donald Fraser.

### F.9. cube

This module implements a data type *cube* for representing multidimensional cubes.

#### F.9.1. Syntax

Table F-3 shows the valid external representations for the *cube* type. *x*, *y*, etc. denote floating-point numbers.

<table>
<thead>
<tr>
<th>External Syntax</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td><em>x</em></td>
<td>A one-dimensional point (or, zero-length one-dimensional interval)</td>
</tr>
<tr>
<td>(x)</td>
<td>Same as above</td>
</tr>
<tr>
<td><em>x1</em>, <em>x2</em>, ..., <em>xn</em></td>
<td>A point in n-dimensional space, represented internally as a zero-volume cube</td>
</tr>
<tr>
<td>(x1, x2, ..., xn)</td>
<td>Same as above</td>
</tr>
<tr>
<td>(x), (y)</td>
<td>A one-dimensional interval starting at <em>x</em> and ending at <em>y</em> or vice versa; the order does not matter</td>
</tr>
<tr>
<td>[(x), (y)]</td>
<td>Same as above</td>
</tr>
<tr>
<td>(<em>x1</em>, ..., <em>xn</em>), (<em>y1</em>, ..., <em>yn</em>)</td>
<td>An n-dimensional cube represented by a pair of its diagonally opposite corners</td>
</tr>
<tr>
<td>[(x1, ..., xn), (yi, ..., yn)]</td>
<td>Same as above</td>
</tr>
</tbody>
</table>

It does not matter which order the opposite corners of a cube are entered in. The *cube* functions auto-
matically swap values if needed to create a uniform “lower left — upper right” internal representation. When the corners coincide, cube stores only one corner along with an “is point” flag to avoid wasting space.

White space is ignored on input, so \([ (x), (y) ]\) is the same as \([ ( x ), ( y ) ]\).

**F.9.2. Precision**

Values are stored internally as 64-bit floating point numbers. This means that numbers with more than about 16 significant digits will be truncated.

**F.9.3. Usage**

Table F-4 shows the operators provided for type `cube`.

*Table F-4. Cube Operators*

<table>
<thead>
<tr>
<th>Operator</th>
<th>Result</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>a = b</code></td>
<td>boolean</td>
<td>The cubes a and b are identical.</td>
</tr>
<tr>
<td><code>a &amp;&amp; b</code></td>
<td>boolean</td>
<td>The cubes a and b overlap.</td>
</tr>
<tr>
<td><code>a @&gt; b</code></td>
<td>boolean</td>
<td>The cube a contains the cube b.</td>
</tr>
<tr>
<td><code>a @$ b</code></td>
<td>boolean</td>
<td>The cube a is contained in the cube b.</td>
</tr>
<tr>
<td><code>a &lt; b</code></td>
<td>boolean</td>
<td>The cube a is less than the cube b.</td>
</tr>
<tr>
<td><code>a &lt;= b</code></td>
<td>boolean</td>
<td>The cube a is less than or equal to the cube b.</td>
</tr>
<tr>
<td><code>a &gt; b</code></td>
<td>boolean</td>
<td>The cube a is greater than the cube b.</td>
</tr>
<tr>
<td><code>a &gt;= b</code></td>
<td>boolean</td>
<td>The cube a is greater than or equal to the cube b.</td>
</tr>
<tr>
<td><code>a &lt;&gt; b</code></td>
<td>boolean</td>
<td>The cube a is not equal to the cube b.</td>
</tr>
<tr>
<td><code>a ~&gt; n</code></td>
<td>float8</td>
<td>Get n-th coordinate of cube (counting from 1).</td>
</tr>
<tr>
<td><code>a ~&gt; n</code></td>
<td>float8</td>
<td>Get n-th coordinate of cube in following way: n = 2 * k - 1 means lower bound of k-th dimension, n = 2 * k means upper bound of k-th dimension. This operator is designed for KNN-GiST support.</td>
</tr>
<tr>
<td><code>a @&lt; b</code></td>
<td>float8</td>
<td>Euclidean distance between a and b.</td>
</tr>
<tr>
<td><code>a @# b</code></td>
<td>float8</td>
<td>Taxicab (L-1 metric) distance between a and b.</td>
</tr>
</tbody>
</table>
Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Operator</th>
<th>Result</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>a &lt;-&gt; b</td>
<td>float8</td>
<td>Chebyshev (L-inf metric) distance between a and b.</td>
</tr>
</tbody>
</table>

(Before PostgreSQL 8.2, the containment operators @> and <@ were respectively called @ and ~. These names are still available, but are deprecated and will eventually be retired. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!)

The scalar ordering operators (<, =>, etc) do not make a lot of sense for any practical purpose but sorting. These operators first compare the first coordinates, and if those are equal, compare the second coordinates, etc. They exist mainly to support the b-tree index operator class for cube, which can be useful for example if you would like a UNIQUE constraint on a cube column.

The cube module also provides a GiST index operator class for cube values. A cube GiST index can be used to search for values using the <=, <&>, @>, and @< operators in WHERE clauses.

In addition, a cube GiST index can be used to find nearest neighbors using the metric operators <=>, <#>, and <~> in ORDER BY clauses. For example, the nearest neighbor of the 3-D point (0.5, 0.5, 0.5) could be found efficiently with:

```
SELECT c FROM test ORDER BY c <=> cube(array[0.5,0.5,0.5]) LIMIT 1;
```

The ~> operator can also be used in this way to efficiently retrieve the first few values sorted by a selected coordinate. For example, to get the first few cubes ordered by the first coordinate (lower left corner) ascending one could use the following query:

```
SELECT c FROM test ORDER BY c ~> 1 LIMIT 5;
```

And to get 2-D cubes ordered by the first coordinate of the upper right corner descending:

```
SELECT c FROM test ORDER BY c ~> 3 DESC LIMIT 5;
```

Table F-5 shows the available functions.

Table F-5. Cube Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Result</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>cube(float8)</td>
<td>cube</td>
<td>Makes a one dimensional cube with both coordinates the same.</td>
<td>cube(1) == '(1)'</td>
</tr>
<tr>
<td>cube(float8, float8)</td>
<td>cube</td>
<td>Makes a one dimensional cube.</td>
<td>cube(1,2) == '(1),(2)'</td>
</tr>
<tr>
<td>cube(float8[])</td>
<td>cube</td>
<td>Makes a zero-volume cube using the coordinates defined by the array.</td>
<td>cube(ARRAY[1,2]) == '(1,2)'</td>
</tr>
</tbody>
</table>
### Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Result</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>cube(float8[], float8[])</code></td>
<td><code>cube</code></td>
<td>Makes a cube with upper right and lower left coordinates as defined by the two arrays, which must be of the same length.</td>
<td><code>cube(ARRAY[1,2], ARRAY[3,4]) == '(1,2),(3,4)'</code></td>
</tr>
<tr>
<td><code>cube(cube, float8)</code></td>
<td><code>cube</code></td>
<td>Makes a new cube by adding a dimension on to an existing cube, with the same values for both endpoints of the new coordinate. This is useful for building cubes piece by piece from calculated values.</td>
<td><code>cube('(1,2),(3,4)'::cube, 5) == '(1,2,5),(3,4,5)'</code></td>
</tr>
<tr>
<td><code>cube(cube, float8, float8)</code></td>
<td><code>cube</code></td>
<td>Makes a new cube by adding a dimension on to an existing cube. This is useful for building cubes piece by piece from calculated values.</td>
<td><code>cube('(1,2),(3,4)'::cube, 5, 6) == '(1,2,5),(3,4,6)'</code></td>
</tr>
<tr>
<td><code>cube_dim(cube)</code></td>
<td><code>integer</code></td>
<td>Returns the number of dimensions of the cube.</td>
<td><code>cube_dim('(1,2),(3,4)') == '2'</code></td>
</tr>
<tr>
<td><code>cube_ll_coord(cube, integer)</code></td>
<td><code>float8</code></td>
<td>Returns the n-th coordinate value for the lower left corner of the cube.</td>
<td><code>cube_ll_coord('(1,2),(3,4)', 2) == '2'</code></td>
</tr>
<tr>
<td><code>cube_ur_coord(cube, integer)</code></td>
<td><code>float8</code></td>
<td>Returns the n-th coordinate value for the upper right corner of the cube.</td>
<td><code>cube_ur_coord('(1,2),(3,4)', 2) == '4'</code></td>
</tr>
<tr>
<td><code>cube_is_point(cube)</code></td>
<td><code>boolean</code></td>
<td>Returns true if the cube is a point, that is, the two defining corners are the same.</td>
<td></td>
</tr>
<tr>
<td><code>cube_distance(cube, cube)</code></td>
<td><code>float8</code></td>
<td>Returns the distance between two cubes. If both cubes are points, this is the normal distance function.</td>
<td></td>
</tr>
</tbody>
</table>
## Appendix F: Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Result</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>cube_subset(cube, integer[])</code></td>
<td>cube</td>
<td>Makes a new cube from an existing cube, using a list of dimension indexes from an array. Can be used to extract the endpoints of a single dimension, or to drop dimensions, or to reorder them as desired.</td>
<td><code>cube_subset(cube('(1,3,5),(6,7,8)'), ARRAY[2]) == '(3),(7)'</code>&lt;br&gt;<code>cube_subset(cube('(1,3,5),(6,7,8)'), ARRAY[3,2,1,1]) == '(5,3,1,1),(8,7,6,6)'</code></td>
</tr>
<tr>
<td><code>cube_union(cube, cube)</code></td>
<td>cube</td>
<td>Produces the union of two cubes.</td>
<td></td>
</tr>
<tr>
<td><code>cube_inter(cube, cube)</code></td>
<td>cube</td>
<td>Produces the intersection of two cubes.</td>
<td></td>
</tr>
</tbody>
</table>
### Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Result</th>
<th>Description</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>cube_enlarge(cube, r double, n integer)</code></td>
<td><code>cube</code></td>
<td>Increases the size of the cube by the specified radius <code>r</code> in at least <code>n</code> dimensions. If the radius is negative the cube is shrunk instead. All defined dimensions are changed by the radius <code>r</code>. Lower-left coordinates are decreased by <code>r</code> and upper-right coordinates are increased by <code>r</code>. If a lower-left coordinate is increased to more than the corresponding upper-right coordinate (this can only happen when <code>r &lt; 0</code>) than both coordinates are set to their average. If <code>n</code> is greater than the number of defined dimensions and the cube is being enlarged (<code>r &gt; 0</code>), then extra dimensions are added to make <code>n</code> altogether; 0 is used as the initial value for the extra coordinates. This function is useful for creating bounding boxes around a point for searching for nearby points.</td>
<td><code>cube_enlarge('(1,2),(3,4)', 0.5, 3) == '(0.5,1.5,-0.5),(3.5,4.5,0.5)'</code></td>
</tr>
</tbody>
</table>

#### F.9.4. Defaults

I believe this union:

```sql
select cube_union('0,5,2),(2,3,1)', '0');
cube_union
-------------------
(0, 0, 0),(2, 5, 2)
(1 row)
```

do not contradict common sense, neither does the intersection

```sql
select cube_inter('0,-1),(1,1)', '(-2),(2)');
cube_inter
```

2365
In all binary operations on differently-dimensioned cubes, I assume the lower-dimensional one to be a Cartesian projection, i.e., having zeroes in place of coordinates omitted in the string representation. The above examples are equivalent to:

```sql
cube_union('(0,5,2),(2,3,1)','(0,0,0),(0,0,0)');
cube_inter('(0,-1),(1,1)','(-2,0),(2,0)');
```

The following containment predicate uses the point syntax, while in fact the second argument is internally represented by a box. This syntax makes it unnecessary to define a separate point type and functions for (box,point) predicates.

```sql
select cube_contains('(0,0),(1,1)', '0.5,0.5');
cube_contains
--------------
t
(1 row)
```

### F.9.5. Notes

For examples of usage, see the regression test `sql/cube.sql`.

To make it harder for people to break things, there is a limit of 100 on the number of dimensions of cubes. This is set in `cubedata.h` if you need something bigger.

### F.9.6. Credits

Original author: Gene Selkov, Jr. <selkovjr@mcs.anl.gov>, Mathematics and Computer Science Division, Argonne National Laboratory.

My thanks are primarily to Prof. Joe Hellerstein (http://db.cs.berkeley.edu/jmh/) for elucidating the gist of the GiST (http://gist.cs.berkeley.edu/), and to his former student Andy Dong for his example written for Illustra. I am also grateful to all Postgres developers, present and past, for enabling myself to create my own world and live undisturbed in it. And I would like to acknowledge my gratitude to Argonne Lab and to the U.S. Department of Energy for the years of faithful support of my database research.

Minor updates to this package were made by Bruno Wolff III <bruno@wolff.to> in August/September of 2002. These include changing the precision from single precision to double precision and adding some new functions.

Additional updates were made by Joshua Reich <josh@root.net> in July 2006. These include `cube(float8[], float8[])` and cleaning up the code to use the V1 call protocol instead of the deprecated V0 protocol.
**F.10. dblink**

dblink is a module that supports connections to other PostgreSQL databases from within a database session.

See also postgres_fdw, which provides roughly the same functionality using a more modern and standards-compliant infrastructure.

**dblink_connect**

**Name**

dblink_connect — opens a persistent connection to a remote database

**Synopsis**

dblink_connect(text connstr) returns text
dblink_connect(text connname, text connstr) returns text

**Description**

dblink_connect() establishes a connection to a remote PostgreSQL database. The server and database to be contacted are identified through a standard libpq connection string. Optionally, a name can be assigned to the connection. Multiple named connections can be open at once, but only one unnamed connection is permitted at a time. The connection will persist until closed or until the database session is ended.

The connection string may also be the name of an existing foreign server. It is recommended to use the foreign-data wrapper dblink_fdw when defining the foreign server. See the example below, as well as CREATE SERVER and CREATE USER MAPPING.

**Arguments**

connname

The name to use for this connection; if omitted, an unnamed connection is opened, replacing any existing unnamed connection.

connstr

libpq-style connection info string, for example hostaddr=127.0.0.1 port=5432 dbname=mydb user=postgres password=mypassword options=-csearch_path=. For details see Section 32.1.1. Alternatively, the name of a foreign server.

**Return Value**

Returns status, which is always OK (since any error causes the function to throw an error instead of returning).
Notes

If untrusted users have access to a database that has not adopted a secure schema usage pattern, begin each session by removing publicly-writable schemas from search_path. One could, for example, add options=-csearch_path= to connstr. This consideration is not specific to dblink; it applies to every interface for executing arbitrary SQL commands.

Only superusers may use dblink_connect to create non-password-authenticated connections. If non-superusers need this capability, use dblink_connect_u instead.

It is unwise to choose connection names that contain equal signs, as this opens a risk of confusion with connection info strings in other dblink functions.

Examples

```sql
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
----------------
    OK
(1 row)

SELECT dblink_connect('myconn', 'dbname=postgres options=-csearch_path=');
----------------
    OK
(1 row)

-- FOREIGN DATA WRAPPER functionality
-- Note: local connection must require password authentication for this to work properly
--     Otherwise, you will receive the following error from dblink_connect():
--     -----------------------------------------------
--     ERROR: password is required
--     DETAIL: Non-superuser cannot connect if the server does not request a password
--     HINT: Target server’s authentication method must be changed.
CREATE SERVER fdtest FOREIGN DATA WRAPPER dblink_fdw OPTIONS (hostaddr '127.0.0.1', dbname 'contrib_regression');
CREATE USER regress_dblink_user WITH PASSWORD 'secret';
CREATE USER MAPPING FOR regress_dblink_user SERVER fdtest OPTIONS (user 'regress_dblink_user', password 'secret');
GRANT USAGE ON FOREIGN SERVER fdtest TO regress_dblink_user;
GRANT SELECT ON TABLE foo TO regress_dblink_user;

\set ORIGINAL_USER :USER
\c - regress_dblink_user
SELECT dblink_connect('myconn', 'fdtest');
----------------
    OK
(1 row)

SELECT * FROM dblink('myconn','SELECT * FROM foo') AS t(a int, b text, c text[]);
 a | b | c
---------
0 | a | (a0,b0,c0)
1 | b | (a1,b1,c1)
2 | c | (a2,b2,c2)```
<p>| | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>3</td>
<td>d</td>
<td>{a3,b3,c3}</td>
</tr>
<tr>
<td>4</td>
<td>e</td>
<td>{a4,b4,c4}</td>
</tr>
<tr>
<td>5</td>
<td>f</td>
<td>{a5,b5,c5}</td>
</tr>
<tr>
<td>6</td>
<td>g</td>
<td>{a6,b6,c6}</td>
</tr>
<tr>
<td>7</td>
<td>h</td>
<td>{a7,b7,c7}</td>
</tr>
<tr>
<td>8</td>
<td>i</td>
<td>{a8,b8,c8}</td>
</tr>
<tr>
<td>9</td>
<td>j</td>
<td>{a9,b9,c9}</td>
</tr>
<tr>
<td>10</td>
<td>k</td>
<td>{a10,b10,c10}</td>
</tr>
</tbody>
</table>

(11 rows)

\c - :ORIGINAL_USER

REVOKE USAGE ON FOREIGN SERVER fdtest FROM regress_dblink_user;
REVOKE SELECT ON TABLE foo FROM regress_dblink_user;
DROP USER MAPPING FOR regress_dblink_user SERVER fdtest;
DROP USER regress_dblink_user;
DROP SERVER fdtest;
dblink_connect_u

Name

dblink_connect_u — opens a persistent connection to a remote database, insecurely

Synopsis

dblink_connect_u(text connstr) returns text
dblink_connect_u(text connname, text connstr) returns text

Description

dblink_connect_u() is identical to dblink_connect(), except that it will allow non-superusers to connect using any authentication method.

If the remote server selects an authentication method that does not involve a password, then impersonation and subsequent escalation of privileges can occur, because the session will appear to have originated from the user as which the local PostgreSQL server runs. Also, even if the remote server does demand a password, it is possible for the password to be supplied from the server environment, such as a ~/.pgpass file belonging to the server’s user. This opens not only a risk of impersonation, but the possibility of exposing a password to an untrustworthy remote server. Therefore, dblink_connect_u() is initially installed with all privileges revoked from PUBLIC, making it un-callable except by superusers. In some situations it may be appropriate to grant EXECUTE permission for dblink_connect_u() to specific users who are considered trustworthy, but this should be done with care. It is also recommended that any ~/.pgpass file belonging to the server’s user not contain any records specifying a wildcard host name.

For further details see dblink_connect().

2370
dblink_disconnect

Name
dblink_disconnect — closes a persistent connection to a remote database

Synopsis
dblink_disconnect() returns text
dblink_disconnect(text connname) returns text

Description
dblink_disconnect() closes a connection previously opened by dblink_connect(). The form with no arguments closes an unnamed connection.

Arguments
connname
The name of a named connection to be closed.

Return Value
Returns status, which is always OK (since any error causes the function to throw an error instead of returning).

Examples
SELECT dblink_disconnect();
dblink_disconnect
-------------------
 OK
 (1 row)

SELECT dblink_disconnect('myconn');
dblink_disconnect
-------------------
 OK
 (1 row)
### dblink

**Name**

dblink — executes a query in a remote database

**Synopsis**

dblink(text connname, text sql [, bool fail_on_error]) returns setof record

dblink(text connstr, text sql [, bool fail_on_error]) returns setof record

dblink(text sql [, bool fail_on_error]) returns setof record

**Description**

dblink executes a query (usually a SELECT, but it can be any SQL statement that returns rows) in a remote database.

When two text arguments are given, the first one is first looked up as a persistent connection’s name; if found, the command is executed on that connection. If not found, the first argument is treated as a connection info string as for dblink_connect, and the indicated connection is made just for the duration of this command.

**Arguments**

- **connname**
  Name of the connection to use; omit this parameter to use the unnamed connection.

- **connstr**
  A connection info string, as previously described for dblink_connect.

- **sql**
  The SQL query that you wish to execute in the remote database, for example select * from foo.

- **fail_on_error**
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function returns no rows.

**Return Value**

The function returns the row(s) produced by the query. Since dblink can be used with any query, it is declared to return record, rather than specifying any particular set of columns. This means that you must specify the expected set of columns in the calling query — otherwise PostgreSQL would not know what to expect. Here is an example:

```
SELECT *
FROM dblink('dbname=mydb options=--csearch_path=',
```
The “alias” part of the FROM clause must specify the column names and types that the function will return. (Specifying column names in an alias is actually standard SQL syntax, but specifying column types is a PostgreSQL extension.) This allows the system to understand what * should expand to, and what proname in the WHERE clause refers to, in advance of trying to execute the function. At run time, an error will be thrown if the actual query result from the remote database does not have the same number of columns shown in the FROM clause. The column names need not match, however, and dblink does not insist on exact type matches either. It will succeed so long as the returned data strings are valid input for the column type declared in the FROM clause.

Notes
A convenient way to use dblink with predetermined queries is to create a view. This allows the column type information to be buried in the view, instead of having to spell it out in every query. For example,

```
CREATE VIEW myremote_pg_proc AS
  SELECT *
  FROM dblink('dbname=postgres options=-csearch_path=',
    'select proname, prosrc from pg_proc')
  AS t1(proname name, prosrc text);
```

```
SELECT * FROM myremote_pg_proc WHERE proname LIKE 'bytea%';
```

Examples
```
SELECT * FROM dblink('dbname=postgres options=-csearch_path=',
  'select proname, prosrc from pg_proc')
AS t1(proname name, prosrc text) WHERE proname LIKE 'bytea%';
```

```
<table>
<thead>
<tr>
<th>proname</th>
<th>prosrc</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteacat</td>
<td>byteacat</td>
</tr>
<tr>
<td>byteaeq</td>
<td>byteaeq</td>
</tr>
<tr>
<td>bytealt</td>
<td>bytealt</td>
</tr>
<tr>
<td>byteale</td>
<td>byteale</td>
</tr>
<tr>
<td>byteagt</td>
<td>byteagt</td>
</tr>
<tr>
<td>byteage</td>
<td>byteage</td>
</tr>
<tr>
<td>byteane</td>
<td>byteane</td>
</tr>
<tr>
<td>byteacmp</td>
<td>byteacmp</td>
</tr>
<tr>
<td>bytealike</td>
<td>bytealike</td>
</tr>
<tr>
<td>byteanlike</td>
<td>byteanlike</td>
</tr>
<tr>
<td>byteain</td>
<td>byteain</td>
</tr>
<tr>
<td>byteaout</td>
<td>byteaout</td>
</tr>
</tbody>
</table>
```

```
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
```

```
dblink_connect
----------------
2373
```
SELECT * FROM dblink('SELECT proname, prosrc FROM pg_proc')
  AS t1(proname name, prosrc text) WHERE proname LIKE 'bytea%';

<table>
<thead>
<tr>
<th>proname</th>
<th>prosrc</th>
</tr>
</thead>
<tbody>
<tr>
<td>byteacat</td>
<td>byteacat</td>
</tr>
<tr>
<td>byteaeq</td>
<td>byteaeq</td>
</tr>
<tr>
<td>bytealt</td>
<td>bytealt</td>
</tr>
<tr>
<td>byteale</td>
<td>byteale</td>
</tr>
<tr>
<td>byteagt</td>
<td>byteagt</td>
</tr>
<tr>
<td>byteage</td>
<td>byteage</td>
</tr>
<tr>
<td>byteane</td>
<td>byteane</td>
</tr>
<tr>
<td>byteacmp</td>
<td>byteacmp</td>
</tr>
<tr>
<td>bytealike</td>
<td>bytealike</td>
</tr>
<tr>
<td>byteanlike</td>
<td>byteanlike</td>
</tr>
<tr>
<td>byteain</td>
<td>byteain</td>
</tr>
<tr>
<td>byteaout</td>
<td>byteaout</td>
</tr>
</tbody>
</table>

(12 rows)

SELECT dblink_connect('myconn', 'dbname=regression options=--csearch_path=');

dblink_connect
----------------
OK

(1 row)

SELECT * FROM dblink('myconn', 'SELECT proname, prosrc FROM pg_proc')
  AS t1(proname name, prosrc text) WHERE proname LIKE 'bytea%';

<table>
<thead>
<tr>
<th>proname</th>
<th>prosrc</th>
</tr>
</thead>
<tbody>
<tr>
<td>bytearecv</td>
<td>bytearecv</td>
</tr>
<tr>
<td>byteasend</td>
<td>byteasend</td>
</tr>
<tr>
<td>byteale</td>
<td>byteale</td>
</tr>
<tr>
<td>byteagt</td>
<td>byteagt</td>
</tr>
<tr>
<td>byteage</td>
<td>byteage</td>
</tr>
<tr>
<td>byteane</td>
<td>byteane</td>
</tr>
<tr>
<td>byteacmp</td>
<td>byteacmp</td>
</tr>
<tr>
<td>bytealike</td>
<td>bytealike</td>
</tr>
<tr>
<td>byteanlike</td>
<td>byteanlike</td>
</tr>
<tr>
<td>byteacat</td>
<td>byteacat</td>
</tr>
<tr>
<td>byteaeq</td>
<td>byteaeq</td>
</tr>
<tr>
<td>bytealt</td>
<td>bytealt</td>
</tr>
<tr>
<td>byteain</td>
<td>byteain</td>
</tr>
<tr>
<td>byteaout</td>
<td>byteaout</td>
</tr>
</tbody>
</table>

(14 rows)
**dblink_exec**

**Name**

dblink_exec — executes a command in a remote database

**Synopsis**

dblink_exec(text connname, text sql [, bool fail_on_error]) returns text  
dblink_exec(text connstr, text sql [, bool fail_on_error]) returns text  
dblink_exec(text sql [, bool fail_on_error]) returns text

**Description**

dblink_exec executes a command (that is, any SQL statement that doesn’t return rows) in a remote database.

When two text arguments are given, the first one is first looked up as a persistent connection’s name; if found, the command is executed on that connection. If not found, the first argument is treated as a connection info string as for dblink_connect, and the indicated connection is made just for the duration of this command.

**Arguments**

connname

Name of the connection to use; omit this parameter to use the unnamed connection.

connstr

A connection info string, as previously described for dblink_connect.

sql

The SQL command that you wish to execute in the remote database, for example `insert into foo values(0,'a',("a0","b0","c0").`

fail_on_error

If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function’s return value is set to ERROR.

**Return Value**

Returns status, either the command’s status string or ERROR.

**Examples**

```
SELECT dblink_connect('dbname=blink_test_standby');
```
```


dblink_connect
----------------
OK
(1 row)

SELECT dblink_exec('insert into foo values(21,"z","{"a0","b0","c0"}");');
  dblink_exec
----------------
INSERT 943366 1
(1 row)

SELECT dblink_connect('myconn', 'dbname=regression');
  dblink_connect
----------------
OK
(1 row)

SELECT dblink_exec('myconn', 'insert into foo values(21,"z","{"a0","b0","c0"}");');
  dblink_exec
----------------
INSERT 6432584 1
(1 row)

SELECT dblink_exec('myconn', 'insert into pg_class values ("foo"),false);
  NOTICE: sql error
  DETAIL:  ERROR:  null value in column "relnamespace" violates not-null constraint

  dblink_exec
----------------
ERROR
(1 row)
```
**dblink_open**

**Name**

`dblink_open` — opens a cursor in a remote database

**Synopsis**

```
`dblink_open(text cursorname, text sql [, bool fail_on_error]) returns text`
`dblink_open(text connname, text cursorname, text sql [, bool fail_on_error]) returns text`
```

**Description**

`dblink_open()` opens a cursor in a remote database. The cursor can subsequently be manipulated with `dblink_fetch()` and `dblink_close()`.

**Arguments**

- `connname`
  
  Name of the connection to use; omit this parameter to use the unnamed connection.

- `cursorname`
  
  The name to assign to this cursor.

- `sql`
  
  The `SELECT` statement that you wish to execute in the remote database, for example `select * from pg_class`.

- `fail_on_error`
  
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function’s return value is set to `ERROR`.

**Return Value**

Returns status, either `OK` or `ERROR`.

**Notes**

Since a cursor can only persist within a transaction, `dblink_open` starts an explicit transaction block (BEGIN) on the remote side, if the remote side was not already within a transaction. This transaction will be closed again when the matching `dblink_close` is executed. Note that if you use `dblink_exec` to change data between `dblink_open` and `dblink_close`, and then an error occurs or you use `dblink_disconnect` before `dblink_close`, your change will be lost because the transaction will be aborted.
Examples

```sql
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
OK
(1 row)

SELECT dblink_open('foo', 'select proname, prosrc from pg_proc');
OK
(1 row)
```
**dblink_fetch**

**Name**

`dblink_fetch` — returns rows from an open cursor in a remote database

**Synopsis**

```sql
dblink_fetch(text cursorname, int howmany [, bool fail_on_error]) returns setof record
dblink_fetch(text connname, text cursorname, int howmany [, bool fail_on_error]) returns setof record
```

**Description**

`dblink_fetch` fetches rows from a cursor previously established by `dblink_open`.

**Arguments**

- **connname**
  Name of the connection to use; omit this parameter to use the unnamed connection.
- **cursorname**
  The name of the cursor to fetch from.
- **howmany**
  The maximum number of rows to retrieve. The next `howmany` rows are fetched, starting at the current cursor position, moving forward. Once the cursor has reached its end, no more rows are produced.
- **fail_on_error**
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function returns no rows.

**Return Value**

The function returns the row(s) fetched from the cursor. To use this function, you will need to specify the expected set of columns, as previously discussed for `dblink`.

**Notes**

On a mismatch between the number of return columns specified in the `FROM` clause, and the actual number of columns returned by the remote cursor, an error will be thrown. In this event, the remote cursor is still advanced by as many rows as it would have been if the error had not occurred. The same is true for any other error occurring in the local query after the remote `FETCH` has been done.
**Examples**

```sql
SELECT dblink_connect('dbname=postgres options=-csearch_path=');
  dblink_connect
----------------
  OK
(1 row)

SELECT dblink_open('foo', 'select proname, prosrc from pg_proc where proname like "bytea%"');
  dblink_open
-------------
  OK
(1 row)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);
  funcname | source
----------+----------
 byteacat | byteacat  
 byteacmp | byteacmp  
 byteaeq  | byteaeq   
 byteage  | byteage   
 byteagt  | byteagt   
(5 rows)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);
  funcname | source
----------+----------
 byteain  | byteain  
 byteale  | byteale  
 bytealike| bytealike 
 bytealt  | bytealt  
 byteane  | byteane  
(5 rows)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);
  funcname | source
----------+----------
 byteanlike| byteanlike
 byteaout  | byteaout 
(2 rows)

SELECT * FROM dblink_fetch('foo', 5) AS (funcname name, source text);
  funcname | source
----------+----------
(0 rows)
```
**dblink_close**

**Name**

`dblink_close` — closes a cursor in a remote database

**Synopsis**

```plaintext
dblink_close(text cursorname [, bool fail_on_error]) returns text
dblink_close(text connname, text cursorname [, bool fail_on_error]) returns text
```

**Description**

`dblink_close` closes a cursor previously opened with `dblink_open`.

**Arguments**

- **connname**
  Name of the connection to use; omit this parameter to use the unnamed connection.
- **cursorname**
  The name of the cursor to close.
- **fail_on_error**
  If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function’s return value is set to `ERROR`.

**Return Value**

Returns status, either `OK` or `ERROR`.

**Notes**

If `dblink_open` started an explicit transaction block, and this is the last remaining open cursor in this connection, `dblink_close` will issue the matching `COMMIT`.

**Examples**

```sql
SELECT dblink_connect('dbname=postgres options=--csearch_path=');
dblink_connect
----------
OK
(1 row)
```
SELECT dblink_open('foo', 'select proname, prosrc from pg_proc');
  dblink_open
------------
  OK
(1 row)

SELECT dblink_close('foo');
  dblink_close
------------
  OK
(1 row)
dblink_get_connections

Name

dblink_get_connections — returns the names of all open named dblink connections

Synopsis

dblink_get_connections() returns text[]

Description

dblink_get_connections returns an array of the names of all open named dblink connections.

Return Value

Returns a text array of connection names, or NULL if none.

Examples

SELECT dblink_get_connections();
### dblink_error_message

#### Name

dblink_error_message — gets last error message on the named connection

#### Synopsis

dblink_error_message(text connname) returns text

#### Description

dblink_error_message fetches the most recent remote error message for a given connection.

#### Arguments

connname

Name of the connection to use.

#### Return Value

Returns last error message, or OK if there has been no error in this connection.

#### Notes

When asynchronous queries are initiated by dblink_send_query, the error message associated with the connection might not get updated until the server’s response message is consumed. This typically means that dblink_is_busy or dblink_get_result should be called prior to dblink_error_message, so that any error generated by the asynchronous query will be visible.

#### Examples

SELECT dblink_error_message(‘dtest1’);
dblink_send_query

Name

dbname_send_query — sends an async query to a remote database

Synopsis

dbtext_send_query(text connname, text sql) returns int

Description

dbtext_send_query sends a query to be executed asynchronously, that is, without immediately waiting for the result. There must not be an async query already in progress on the connection.

After successfully dispatching an async query, completion status can be checked with dbtext_is_busy, and the results are ultimately collected with dbtext_get_result. It is also possible to attempt to cancel an active async query using dbtext_cancel_query.

Arguments

connname

Name of the connection to use.

sql

The SQL statement that you wish to execute in the remote database, for example select * from pg_class.

Return Value

Returns 1 if the query was successfully dispatched, 0 otherwise.

Examples

SELECT dbtext_send_query('dtest1', 'SELECT * FROM foo WHERE f1 < 3');


dblink_is_busy

Name
dblink_is_busy — checks if connection is busy with an async query

Synopsis
dblink_is_busy(text connname) returns int

Description
dblink_is_busy tests whether an async query is in progress.

Arguments

connname

Name of the connection to check.

Return Value

Returns 1 if connection is busy, 0 if it is not busy. If this function returns 0, it is guaranteed that
dblink_get_result will not block.

Examples

SELECT dblink_is_busy('dtest1');
**dblink_get_notify**

**Name**

`dblink_get_notify` — retrieve async notifications on a connection

**Synopsis**

```
dblink_get_notify() returns setof (notify_name text, be_pid int, extra text)
dblink_get_notify(text connname) returns setof (notify_name text, be_pid int, extra text)
```

**Description**

`dblink_get_notify` retrieves notifications on either the unnamed connection, or on a named connection if specified. To receive notifications via `dblink`, `LISTEN` must first be issued, using `dblink_exec`. For details see `LISTEN` and `NOTIFY`.

**Arguments**

`connname`

The name of a named connection to get notifications on.

**Return Value**

Returns `setof (notify_name text, be_pid int, extra text)`, or an empty set if none.

**Examples**

```sql
SELECT dblink_exec('LISTEN virtual');
dblink_exec
-------------
LISTEN
(1 row)

SELECT * FROM dblink_get_notify();
notify_name | be_pid | extra
-------------+--------+-------
(0 rows)

NOTIFY virtual;
NOTIFY

SELECT * FROM dblink_get_notify();
notify_name | be_pid | extra
-------------+--------+-------
virtual | 1229 |
(1 row)
```
**dblink_get_result**

**Name**

dblink_get_result — gets an async query result

**Synopsis**

dblink_get_result(text connname [, bool fail_on_error]) returns setof record

**Description**

dblink_get_result collects the results of an asynchronous query previously sent with dblink_send_query. If the query is not already completed, dblink_get_result will wait until it is.

**Arguments**

connname

Name of the connection to use.

fail_on_error

If true (the default when omitted) then an error thrown on the remote side of the connection causes an error to also be thrown locally. If false, the remote error is locally reported as a NOTICE, and the function returns no rows.

**Return Value**

For an async query (that is, a SQL statement returning rows), the function returns the row(s) produced by the query. To use this function, you will need to specify the expected set of columns, as previously discussed for dblink.

For an async command (that is, a SQL statement not returning rows), the function returns a single row with a single text column containing the command’s status string. It is still necessary to specify that the result will have a single text column in the calling FROM clause.

**Notes**

This function must be called if dblink_send_query returned 1. It must be called once for each query sent, and one additional time to obtain an empty set result, before the connection can be used again.

When using dblink_send_query and dblink_get_result, dblink fetches the entire remote query result before returning any of it to the local query processor. If the query returns a large number of rows, this can result in transient memory bloat in the local session. It may be better to open such
a query as a cursor with `dblink_open` and then fetch a manageable number of rows at a time. Alternatively, use plain `dblink()`, which avoids memory bloat by spooling large result sets to disk.

**Examples**

```sql
contrib_regression=# SELECT dblink_connect('dtest1', 'dbname=contrib_regression');
  dblink_connect
----------------
    OK
(1 row)

contrib_regression=# SELECT * FROM contrib_regression-
contrib_regression=# dblink_send_query('dtest1', 'select * from foo where f1 < 3') AS t1;
t1    ----
      1
(1 row)

contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS t1(f1 int, f2 text, f3 text[
  f1 | f2 | f3
----+----+------------
  0 | a  | {a0,b0,c0}
  1 | b  | {a1,b1,c1}
  2 | c  | {a2,b2,c2}
(3 rows)

contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS t1(f1 int, f2 text, f3 text[
  f1 | f2 | f3
----+----+---------------
  7 | h  | {a7,b7,c7}
  8 | i  | {a8,b8,c8}
  9 | j  | {a9,b9,c9}
10 | k  | {a10,b10,c10}
(4 rows)
```
```sql
contrib_regression=# SELECT * FROM dblink_get_result('dtest1') AS t1(f1 int, f2 text, f3 text)
  f1 | f2 | f3
----+----+----
(0 rows)
```
The function `dblink_cancel_query` is used to cancel any active query on the named connection. Its synopsis is:

```sql
dblink_cancel_query(text connname) returns text
```

**Description**

`dblink_cancel_query` attempts to cancel any query that is in progress on the named connection. Note that this is not certain to succeed (since, for example, the remote query might already have finished). A cancel request simply improves the odds that the query will fail soon. You must still complete the normal query protocol, for example by calling `dblink_get_result`.

**Arguments**

- **connname**
  - Name of the connection to use.

**Return Value**

Returns `OK` if the cancel request has been sent, or the text of an error message on failure.

**Examples**

```sql
SELECT dblink_cancel_query('dtest1');
```
**dblink_get_pkey**

**Name**

dblink_get_pkey — returns the positions and field names of a relation’s primary key fields

**Synopsis**

dblink_get_pkey(text relname) returns setof dblink_pkey_results

**Description**

dblink_get_pkey provides information about the primary key of a relation in the local database. This is sometimes useful in generating queries to be sent to remote databases.

**Arguments**

rename

Name of a local relation, for example foo or myschema.mytab. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

**Return Value**

Returns one row for each primary key field, or no rows if the relation has no primary key. The result row type is defined as

CREATE TYPE dblink_pkey_results AS (position int, colname text);

The position column simply runs from 1 to N; it is the number of the field within the primary key, not the number within the table’s columns.

**Examples**

CREATE TABLE foobar {
   f1 int,
   f2 int,
   f3 int,
   PRIMARY KEY (f1, f2, f3)
};
CREATE TABLE

SELECT * FROM dblink_get_pkey('foobar');

<table>
<thead>
<tr>
<th>position</th>
<th>colname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>f1</td>
</tr>
<tr>
<td>2</td>
<td>f2</td>
</tr>
</tbody>
</table>
### dblink_get_pkey

<table>
<thead>
<tr>
<th></th>
<th>f3</th>
</tr>
</thead>
</table>

(3 rows)
**dblink_build_sql_insert**

**Name**

`dblink_build_sql_insert` — builds an INSERT statement using a local tuple, replacing the primary key field values with alternative supplied values.

**Synopsis**

```sql

dblink_build_sql_insert(text relname,
    int2vector primary_key_attnums,
    integer num_primary_key_atts,
    text[] src_pk_att_vals_array,
    text[] tgt_pk_att_vals_array) returns text
```

**Description**

`dblink_build_sql_insert` can be useful in doing selective replication of a local table to a remote database. It selects a row from the local table based on primary key, and then builds a SQL **INSERT** command that will duplicate that row, but with the primary key values replaced by the values in the last argument. (To make an exact copy of the row, just specify the same values for the last two arguments.)

**Arguments**

- `relname`
  
  Name of a local relation, for example `foo` or `myschema.mytab`. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

- `primary_key_attnums`
  
  Attribute numbers (1-based) of the primary key fields, for example `1` `2`.

- `num_primary_key_atts`
  
  The number of primary key fields.

- `src_pk_att_vals_array`
  
  Values of the primary key fields to be used to look up the local tuple. Each field is represented in text form. An error is thrown if there is no local row with these primary key values.

- `tgt_pk_att_vals_array`
  
  Values of the primary key fields to be placed in the resulting **INSERT** command. Each field is represented in text form.

**Return Value**

Returns the requested SQL statement as text.
Notes

As of PostgreSQL 9.0, the attribute numbers in primary_key_attnums are interpreted as logical column numbers, corresponding to the column’s position in SELECT * FROM relname. Previous versions interpreted the numbers as physical column positions. There is a difference if any column(s) to the left of the indicated column have been dropped during the lifetime of the table.

Examples

```sql
SELECT dblink_build_sql_insert('foo', '1 2', 2, '{"1", "a"}', '{"1", "b""a"}');
```

```
INSERT INTO foo(f1,f2,f3) VALUES('1','b"a','1')
(1 row)
```
dblink_build_sql_delete

Name
dblink_build_sql_delete — builds a DELETE statement using supplied values for primary key field values

Synopsis
dblink_build_sql_delete(text relname,
                        int2vector primary_key_attnums,
                        integer num_primary_key_atts,
                        text[] tgt_pk_att_vals_array) returns text

Description
dblink_build_sql_delete can be useful in doing selective replication of a local table to a remote database. It builds a SQL DELETE command that will delete the row with the given primary key values.

Arguments

rename
    Name of a local relation, for example foo or myschema.mytab. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

primary_key_attnums
    Attribute numbers (1-based) of the primary key fields, for example 1 2.
	num_primary_key_atts
    The number of primary key fields.

tgt_pk_att_vals_array
    Values of the primary key fields to be used in the resulting DELETE command. Each field is represented in text form.

Return Value
Returns the requested SQL statement as text.

Notes
As of PostgreSQL 9.0, the attribute numbers in primary_key_attnums are interpreted as logical column numbers, corresponding to the column’s position in SELECT * FROM relname. Previous
versions interpreted the numbers as physical column positions. There is a difference if any column(s) to the left of the indicated column have been dropped during the lifetime of the table.

Examples

```sql
SELECT dblink_build_sql_delete('"MyFoo"', '1 2', 2, '{"1", "b"}');
```

```sql
DELETE FROM "MyFoo" WHERE f1='1' AND f2='b'
(1 row)
```
**dblink_build_sql_update**

**Name**

`dblink_build_sql_update` — builds an UPDATE statement using a local tuple, replacing the primary key field values with alternative supplied values.

**Synopsis**

```sql
dblink_build_sql_update(text relname,
        int2vector primary_key_attnums,
        integer num_primary_key_atts,
        text[] src_pk_att_vals_array,
        text[] tgt_pk_att vals_array) returns text
```

**Description**

`dblink_build_sql_update` can be useful in doing selective replication of a local table to a remote database. It selects a row from the local table based on primary key, and then builds a SQL UPDATE command that will duplicate that row, but with the primary key values replaced by the values in the last argument. (To make an exact copy of the row, just specify the same values for the last two arguments.) The UPDATE command always assigns all fields of the row — the main difference between this and `dblink_build_sql_insert` is that it’s assumed that the target row already exists in the remote table.

**Arguments**

- **relname**
  
  Name of a local relation, for example `foo` or `myschema.mytab`. Include double quotes if the name is mixed-case or contains special characters, for example "FooBar"; without quotes, the string will be folded to lower case.

- **primary_key_attnums**
  
  Attribute numbers (1-based) of the primary key fields, for example 1 2.

- **num_primary_key_atts**
  
  The number of primary key fields.

- **src_pk_att_vals_array**
  
  Values of the primary key fields to be used to look up the local tuple. Each field is represented in text form. An error is thrown if there is no local row with these primary key values.

- **tgt_pk_att_vals_array**
  
  Values of the primary key fields to be placed in the resulting UPDATE command. Each field is represented in text form.
Return Value

Returns the requested SQL statement as text.

Notes

As of PostgreSQL 9.0, the attribute numbers in primary_key_attnums are interpreted as logical column numbers, corresponding to the column’s position in SELECT * FROM relname. Previous versions interpreted the numbers as physical column positions. There is a difference if any column(s) to the left of the indicated column have been dropped during the lifetime of the table.

Examples

```sql
SELECT dblink_build_sql_update('foo', '1 2', 2, '{"1", "a"}', '{"1", "b"}');
```

```
UPDATE foo SET f1='1', f2='b', f3='1' WHERE f1='1' AND f2='b' (1 row)
```
F.11. dict_int

dict_int is an example of an add-on dictionary template for full-text search. The motivation for this example dictionary is to control the indexing of integers (signed and unsigned), allowing such numbers to be indexed while preventing excessive growth in the number of unique words, which greatly affects the performance of searching.

F.11.1. Configuration

The dictionary accepts two options:

- The maxlen parameter specifies the maximum number of digits allowed in an integer word. The default value is 6.
- The rejectlong parameter specifies whether an overlength integer should be truncated or ignored. If rejectlong is false (the default), the dictionary returns the first maxlen digits of the integer. If rejectlong is true, the dictionary treats an overlength integer as a stop word, so that it will not be indexed. Note that this also means that such an integer cannot be searched for.

F.11.2. Usage

Installing the dict_int extension creates a text search template intdict_template and a dictionary intdict based on it, with the default parameters. You can alter the parameters, for example

```sql
mydb# ALTER TEXT SEARCH DICTIONARY intdict (MAXLEN = 4, REJECTLONG = true);
```

or create new dictionaries based on the template.

To test the dictionary, you can try

```sql
mydb# select ts_lexize('intdict', '12345678');
```

but real-world usage will involve including it in a text search configuration as described in Chapter 12. That might look like this:

```sql
ALTER TEXT SEARCH CONFIGURATION english
  ALTER MAPPING FOR int, uint WITH intdict;
```

F.12. dict_xsyn

dict_xsyn (Extended Synonym Dictionary) is an example of an add-on dictionary template for full-text search. This dictionary type replaces words with groups of their synonyms, and so makes it possible to search for a word using any of its synonyms.
Appendix F. Additional Supplied Modules

F.12.1. Configuration

A `dict_xsyn` dictionary accepts the following options:

- `matchorig` controls whether the original word is accepted by the dictionary. Default is `true`.
- `matchsynonyms` controls whether the synonyms are accepted by the dictionary. Default is `false`.
- `keeporig` controls whether the original word is included in the dictionary's output. Default is `true`.
- `keepsynonyms` controls whether the synonyms are included in the dictionary's output. Default is `true`.
- `rules` is the base name of the file containing the list of synonyms. This file must be stored in `$SHAREDIR/tsearch_data/` (where `$SHAREDIR` means the PostgreSQL installation’s shared-data directory). Its name must end in `.rules` (which is not to be included in the `rules` parameter).

The rules file has the following format:

- Each line represents a group of synonyms for a single word, which is given first on the line. Synonyms are separated by whitespace, thus:
  ```
  word syn1 syn2 syn3
  ```
- The sharp (`#`) sign is a comment delimiter. It may appear at any position in a line. The rest of the line will be skipped.

Look at `xsyn_sample.rules`, which is installed in `$SHAREDIR/tsearch_data/`, for an example.

F.12.2. Usage

Installing the `dict_xsyn` extension creates a text search template `xsyn_template` and a dictionary `xsyn` based on it, with default parameters. You can alter the parameters, for example

```sql
mydb# ALTER TEXT SEARCH DICTIONARY xsyn (RULES='my_rules', KEEPORIG=false);
ALTER TEXT SEARCH DICTIONARY
```

or create new dictionaries based on the template.

To test the dictionary, you can try

```sql
mydb=# SELECT ts_lexize('xsyn', 'word');
```

```
{syn1,syn2,syn3}
```

```sql
mydb# ALTER TEXT SEARCH DICTIONARY xsyn (RULES='my_rules', KEEPORIG=true);
ALTER TEXT SEARCH DICTIONARY
```

```sql
mydb=# SELECT ts_lexize('xsyn', 'word');
```

```
{word,syn1,syn2,syn3}
```

```sql
mydb# ALTER TEXT SEARCH DICTIONARY xsyn (RULES='my_rules', KEEPORIG=false, MATCHSYNONYMS=true);
ALTER TEXT SEARCH DICTIONARY
```

```sql
mydb=# SELECT ts_lexize('xsyn', 'syn1');
```

...
Appendix F. Additional Supplied Modules

Real-world usage will involve including it in a text search configuration as described in Chapter 12. That might look like this:

```sql
ALTER TEXT SEARCH CONFIGURATION english
  ALTER MAPPING FOR word, asciword WITH xsyn, english_stem;
```

## F.13. earthdistance

The `earthdistance` module provides two different approaches to calculating great circle distances on the surface of the Earth. The one described first depends on the `cube` module (which must be installed before `earthdistance` can be installed). The second one is based on the built-in `point` data type, using longitude and latitude for the coordinates.

In this module, the Earth is assumed to be perfectly spherical. (If that’s too inaccurate for you, you might want to look at the PostGIS project.)

### F.13.1. Cube-based Earth Distances

Data is stored in cubes that are points (both corners are the same) using 3 coordinates representing the x, y, and z distance from the center of the Earth. A domain `earth over cube` is provided, which includes constraint checks that the value meets these restrictions and is reasonably close to the actual surface of the Earth.

The radius of the Earth is obtained from the `earth()` function. It is given in meters. But by changing this one function you can change the module to use some other units, or to use a different value of the radius that you feel is more appropriate.

This package has applications to astronomical databases as well. Astronomers will probably want to change `earth()` to return a radius of 180/\(\pi\) so that distances are in degrees.

Functions are provided to support input in latitude and longitude (in degrees), to support output of latitude and longitude, to calculate the great circle distance between two points and to easily specify a bounding box usable for index searches.

The provided functions are shown in Table F-6.

### Table F-6. Cube-based Earthdistance Functions

2. [http://postgis.net/](http://postgis.net/)
Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>earth()</td>
<td>float8</td>
<td>Returns the assumed radius of the Earth.</td>
</tr>
<tr>
<td>sec_to_gc(float8)</td>
<td>float8</td>
<td>Converts the normal straight line (secant) distance between two points on the surface of the Earth to the great circle distance between them.</td>
</tr>
<tr>
<td>gc_to_sec(float8)</td>
<td>float8</td>
<td>Converts the great circle distance between two points on the surface of the Earth to the normal straight line (secant) distance between them.</td>
</tr>
<tr>
<td>ll_to_earth(float8, float8)</td>
<td>earth</td>
<td>Returns the location of a point on the surface of the Earth given its latitude (argument 1) and longitude (argument 2) in degrees.</td>
</tr>
<tr>
<td>latitude(earth)</td>
<td>float8</td>
<td>Returns the latitude in degrees of a point on the surface of the Earth.</td>
</tr>
<tr>
<td>longitude(earth)</td>
<td>float8</td>
<td>Returns the longitude in degrees of a point on the surface of the Earth.</td>
</tr>
<tr>
<td>earth_distance(earth, earth)</td>
<td>float8</td>
<td>Returns the great circle distance between two points on the surface of the Earth.</td>
</tr>
<tr>
<td>earth_box(earth, float8)</td>
<td>cube</td>
<td>Returns a box suitable for an indexed search using the cube $@ &gt;$ operator for points within a given great circle distance of a location. Some points in this box are further than the specified great circle distance from the location, so a second check using earth_distance should be included in the query.</td>
</tr>
</tbody>
</table>

F.13.2. Point-based Earth Distances

The second part of the module relies on representing Earth locations as values of type point, in which the first component is taken to represent longitude in degrees, and the second component is taken to represent latitude in degrees. Points are taken as (longitude, latitude) and not vice versa because longitude is closer to the intuitive idea of x-axis and latitude to y-axis.

A single operator is provided, shown in Table F-7.
Table F-7. Point-based Earthdistance Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>point &lt;@&gt; point</td>
<td>float8</td>
<td>Gives the distance in statute miles between two points on the Earth’s surface.</td>
</tr>
</tbody>
</table>

Note that unlike the cube-based part of the module, units are hardwired here: changing the `earth()` function will not affect the results of this operator.

One disadvantage of the longitude/latitude representation is that you need to be careful about the edge conditions near the poles and near +/- 180 degrees of longitude. The cube-based representation avoids these discontinuities.

**F.14. file_fdw**

The `file_fdw` module provides the foreign-data wrapper `file_fdw`, which can be used to access data files in the server’s file system. Data files must be in a format that can be read by `COPY FROM`; see `COPY` for details. Access to such data files is currently read-only.

A foreign table created using this wrapper can have the following options:

- **filename**
  - Specifies the file to be read. Required. Must be an absolute path name.

- **format**
  - Specifies the file’s format, the same as `COPY`’s `FORMAT` option.

- **header**
  - Specifies whether the file has a header line, the same as `COPY`’s `HEADER` option.

- **delimiter**
  - Specifies the file’s delimiter character, the same as `COPY`’s `DELIMITER` option.

- **quote**
  - Specifies the file’s quote character, the same as `COPY`’s `QUOTE` option.

- **escape**
  - Specifies the file’s escape character, the same as `COPY`’s `ESCAPE` option.

- **null**
  - Specifies the file’s null string, the same as `COPY`’s `NULL` option.

- **encoding**
  - Specifies the file’s encoding, the same as `COPY`’s `ENCODING` option.

Note that while `COPY` allows options such as OIDS and HEADER to be specified without a corresponding value, the foreign data wrapper syntax requires a value to be present in all cases. To activate `COPY` options normally supplied without a value, you can instead pass the value `TRUE`.

A column of a foreign table created using this wrapper can have the following options:
force_not_null

This is a Boolean option. If true, it specifies that values of the column should not be matched against the null string (that is, the file-level null option). This has the same effect as listing the column in COPY’s FORCE_NOT_NULL option.

force_null

This is a Boolean option. If true, it specifies that values of the column which match the null string are returned as NULL even if the value is quoted. Without this option, only unquoted values matching the null string are returned as NULL. This has the same effect as listing the column in COPY’s FORCE_NULL option.

COPY’s OIDS and FORCE_QUOTE options are currently not supported by file_fdw.

These options can only be specified for a foreign table or its columns, not in the options of the file_fdw foreign-data wrapper, nor in the options of a server or user mapping using the wrapper.

Changing table-level options requires superuser privileges, for security reasons: only a superuser should be able to determine which file is read. In principle non-superusers could be allowed to change the other options, but that’s not supported at present.

For a foreign table using file_fdw, EXPLAIN shows the name of the file to be read. Unless COSTS OFF is specified, the file size (in bytes) is shown as well.

Example F-1. Create a Foreign Table for PostgreSQL CSV Logs

One of the obvious uses for the file_fdw is to make the PostgreSQL activity log available as a table for querying. To do this, first you must be logging to a CSV file, which here we will call pglog.csv.

First, install file_fdw as an extension:

```
CREATE EXTENSION file_fdw;
```

Then create a foreign server:

```
CREATE SERVER pglog FOREIGN DATA WRAPPER file_fdw;
```

Now you are ready to create the foreign data table. Using the CREATE FOREIGN TABLE command, you will need to define the columns for the table, the CSV file name, and its format:

```
CREATE FOREIGN TABLE pglog (  
  log_time timestamp(3) with time zone,  
  user_name text,  
  database_name text,  
  process_id integer,  
  connection_from text,  
  session_id text,  
  session_line_num bigint,  
  command_tag text,  
  session_start_time timestamp with time zone,  
  virtual_transaction_id text,  
  transaction_id bigint,  
  error_severity text,  
  sql_state_code text,  
  message text,  
  detail text,  
  hint text,  
  internal_query text,  
  internal_query_pos integer,  
  context text,  
  query text,  
  query_pos integer,
```
location text,
  application_name text
) SERVER pglog
OPTIONS ( filename '/home/josh/9.1/data/pg_log/pglog.csv', format 'csv' );

That’s it — now you can query your log directly. In production, of course, you would need to define some way to deal with log rotation.

F.15. fuzzystrmatch

The fuzzystrmatch module provides several functions to determine similarities and distance between strings.

Caution
At present, the soundex, metaphone, dmetaphone, and dmetaphone_alt functions do not work well with multibyte encodings (such as UTF-8).

F.15.1. Soundex

The Soundex system is a method of matching similar-sounding names by converting them to the same code. It was initially used by the United States Census in 1880, 1900, and 1910. Note that Soundex is not very useful for non-English names.

The fuzzystrmatch module provides two functions for working with Soundex codes:

soundex(text) returns text
difference(text, text) returns int

The soundex function converts a string to its Soundex code. The difference function converts two strings to their Soundex codes and then reports the number of matching code positions. Since Soundex codes have four characters, the result ranges from zero to four, with zero being no match and four being an exact match. (Thus, the function is misnamed — similarity would have been a better name.)

Here are some usage examples:

SELECT soundex('hello world!');

SELECT soundex('Anne'), soundex('Ann'), difference('Anne', 'Ann');
SELECT soundex('Anne'), soundex('Andrew'), difference('Anne', 'Andrew');
SELECT soundex('Anne'), soundex('Margaret'), difference('Anne', 'Margaret');

CREATE TABLE s (nm text);
INSERT INTO s VALUES ('john');
INSERT INTO s VALUES ('joan');
INSERT INTO s VALUES ('wobbly');
INSERT INTO s VALUES ('jack');

SELECT * FROM s WHERE soundex(nm) = soundex('john');
F.15.2. Levenshtein

This function calculates the Levenshtein distance between two strings:

\[
\text{levenshtein}(\text{source}, \text{target}, \text{ins\_cost}, \text{del\_cost}, \text{sub\_cost}) \quad \text{returns int}
\]

\[
\text{levenshtein}(\text{source}, \text{target}) \quad \text{returns int}
\]

\[
\text{levenshtein\_less\_equal}(\text{source}, \text{target}, \text{ins\_cost}, \text{del\_cost}, \text{sub\_cost}, \text{max\_d}) \quad \text{returns int}
\]

\[
\text{levenshtein\_less\_equal}(\text{source}, \text{target}, \text{max\_d}) \quad \text{returns int}
\]

Both source and target can be any non-null string, with a maximum of 255 characters. The cost parameters specify how much to charge for a character insertion, deletion, or substitution, respectively. You can omit the cost parameters, as in the second version of the function; in that case they all default to 1.

\text{levenshtein\_less\_equal} is an accelerated version of the Levenshtein function for use when only small distances are of interest. If the actual distance is less than or equal to \text{max\_d}, then \text{levenshtein\_less\_equal} returns the correct distance; otherwise it returns some value greater than \text{max\_d}. If \text{max\_d} is negative then the behavior is the same as \text{levenshtein}.

Examples:

\[
\text{test=} # \quad \text{SELECT levenshtein(\text{'GUMBO'}, \text{'GAMBOL'});}
\]

\[
\text{levenshtein} \quad \text{------------------------}
\]

\[
\text{2}
\]

(1 row)

\[
\text{test=} # \quad \text{SELECT levenshtein(\text{'GUMBO'}, \text{'GAMBOL'}, 2,1,1);}
\]

\[
\text{levenshtein} \quad \text{------------------------}
\]

\[
\text{3}
\]

(1 row)

\[
\text{test=} # \quad \text{SELECT levenshtein\_less\_equal(\text{'extensive'}, \text{'exhaustive',2});}
\]

\[
\text{levenshtein\_less\_equal} \quad \text{------------------------}
\]

\[
\text{3}
\]

(1 row)

\[
\text{test=} # \quad \text{SELECT levenshtein\_less\_equal(\text{'extensive'}, \text{'exhaustive',4});}
\]

\[
\text{levenshtein\_less\_equal} \quad \text{------------------------}
\]

\[
\text{4}
\]

(1 row)

F.15.3. Metaphone

Metaphone, like Soundex, is based on the idea of constructing a representative code for an input string. Two strings are then deemed similar if they have the same codes.

This function calculates the metaphone code of an input string:
metaphone(text source, int max_output_length) returns text

source has to be a non-null string with a maximum of 255 characters. max_output_length sets the maximum length of the output metaphone code; if longer, the output is truncated to this length.

Example:

test=# SELECT metaphone('GUMBO', 4);
metaphone
-----------
KM
(1 row)

F.15.4. Double Metaphone

The Double Metaphone system computes two “sounds like” strings for a given input string — a “primary” and an “alternate”. In most cases they are the same, but for non-English names especially they can be a bit different, depending on pronunciation. These functions compute the primary and alternate codes:

dmetaphone(text source) returns text
dmetaphone_alt(text source) returns text

There is no length limit on the input strings.

Example:

test=# select dmetaphone('gumbo');
dmetaphone
------------
KMP
(1 row)

F.16. hstore

This module implements the hstore data type for storing sets of key/value pairs within a single PostgreSQL value. This can be useful in various scenarios, such as rows with many attributes that are rarely examined, or semi-structured data. Keys and values are simply text strings.

F.16.1. hstore External Representation

The text representation of an hstore, used for input and output, includes zero or more key => value pairs separated by commas. Some examples:

k => v
foo => bar, baz => whatever
"1-a" => "anything at all"

The order of the pairs is not significant (and may not be reproduced on output). Whitespace between pairs or around the => sign is ignored. Double-quote keys and values that include whitespace, commas, -s or >s. To include a double quote or a slash in a key or value, escape it with a backslash.
Each key in an hstore is unique. If you declare an hstore with duplicate keys, only one will be stored in the hstore and there is no guarantee as to which will be kept:

```sql
SELECT 'a=>1,a=>2'::hstore;

hstore
----------------
"a"=>"1"
```

A value (but not a key) can be an SQL NULL. For example:

key => NULL

The NULL keyword is case-insensitive. Double-quote the NULL to treat it as the ordinary string “NULL”.

Note: Keep in mind that the hstore text format, when used for input, applies before any required quoting or escaping. If you are passing an hstore literal via a parameter, then no additional processing is needed. But if you’re passing it as a quoted literal constant, then any single-quote characters and (depending on the setting of the standard_conforming_strings configuration parameter) backslash characters need to be escaped correctly. See Section 4.1.2.1 for more on the handling of string constants.

On output, double quotes always surround keys and values, even when it’s not strictly necessary.

### F.16.2. hstore Operators and Functions

The operators provided by the hstore module are shown in Table F-8, the functions in Table F-9.

#### Table F-8. hstore Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>hstore -&gt; text</td>
<td>get value for key (NULL if not present)</td>
<td>‘a=&gt;x, b=&gt;y’::hstore -&gt; ‘a’</td>
<td>x</td>
</tr>
<tr>
<td>hstore -&gt; text[]</td>
<td>get values for keys (NULL if not present)</td>
<td>‘a=&gt;x, b=&gt;y, c=&gt;z’::hstore -&gt; ARRAY[‘c’,”a”]</td>
<td>{&quot;z&quot;,&quot;x&quot;}</td>
</tr>
<tr>
<td>hstore</td>
<td></td>
<td>hstore</td>
<td>concatenate hstores</td>
</tr>
<tr>
<td>hstore ? text</td>
<td>does hstore contain key?</td>
<td>‘a=&gt;1’::hstore ? ‘a’</td>
<td>t</td>
</tr>
<tr>
<td>hstore ?&amp; text[]</td>
<td>does hstore contain all specified keys?</td>
<td>‘a=&gt;1,b=&gt;2’::hstore ?&amp; ARRAY[‘a’,’b’]</td>
<td></td>
</tr>
<tr>
<td>hstore ?</td>
<td>text[]</td>
<td>does hstore contain any of the specified keys?</td>
<td>‘a=&gt;1,b=&gt;2’::hstore ?</td>
</tr>
</tbody>
</table>
### Appendix F. Additional Supplied Modules

#### Table F-9. hstore Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>hstore(record)</td>
<td>hstore</td>
<td>construct an hstore from a record or row</td>
<td>hstore(ROW(1,2) f1 =&gt; 1, f2 =&gt; 2)</td>
<td></td>
</tr>
<tr>
<td>hstore(text[])</td>
<td>hstore</td>
<td>construct an hstore from an array, which may be either a key/value array, or a two-dimensional array</td>
<td>hstore(ARRAY(['a','b','c','d'])</td>
<td></td>
</tr>
</tbody>
</table>

Note: Prior to PostgreSQL 8.2, the containment operators @> and @@ were called @ and ~, respectively. These names are still available, but are deprecated and will eventually be removed. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>hstore(text[],</td>
<td>hstore</td>
<td>construct an hstore from separate key and value arrays</td>
<td>hstore(Array[('a', 1), ('b', 2)])</td>
<td>&quot;a&quot;=&gt;&quot;1&quot;, &quot;b&quot;=&gt;&quot;2&quot;</td>
</tr>
<tr>
<td>text[])</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>hstore(text,</td>
<td>hstore</td>
<td>make single-item hstore</td>
<td>hstore('a', 'b')</td>
<td>&quot;a&quot;=&gt;&quot;b&quot;</td>
</tr>
<tr>
<td>text)</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>akeys(hstore)</td>
<td>text[]</td>
<td>get hstore's keys as an array</td>
<td>akeys('a=&gt;1, b=&gt;2')</td>
<td>{a, b}</td>
</tr>
<tr>
<td>skeys(hstore)</td>
<td>setof text</td>
<td>get hstore's keys as a set</td>
<td>skeys('a=&gt;1, b=&gt;2')</td>
<td>a b</td>
</tr>
<tr>
<td>avals(hstore)</td>
<td>text[]</td>
<td>get hstore's values as an array</td>
<td>avals('a=&gt;1, b=&gt;2')</td>
<td>{1, 2}</td>
</tr>
<tr>
<td>svals(hstore)</td>
<td>setof text</td>
<td>get hstore's values as a set</td>
<td>svals('a=&gt;1, b=&gt;2')</td>
<td>1 2</td>
</tr>
<tr>
<td>hstore_to_array</td>
<td>text[]</td>
<td>get hstore's keys and values as an array of alternating keys and values</td>
<td>hstore_to_array(#{a=&gt;1, b=&gt;2})</td>
<td>{a, 1, b, 2}</td>
</tr>
<tr>
<td>hstore_to_matrix</td>
<td>text[]</td>
<td>get hstore's keys and values as a two-dimensional array</td>
<td>hstore_to_matrix(#{a=&gt;1, b=&gt;2}, 2)</td>
<td>x(#{a=&gt;1, b=&gt;2})</td>
</tr>
<tr>
<td>hstore_to_json</td>
<td>json</td>
<td>get hstore as a json value, converting all non-null values to JSON strings</td>
<td>hstore_to_json(&quot;a key&quot;=&gt;1, &quot;b&quot;=&gt;t, &quot;c&quot;=&gt;null, &quot;d&quot;=&gt;12345, &quot;e&quot;=&gt;012345, &quot;f&quot;=&gt;1.234, &quot;g&quot;=&gt;2.345e+4&quot;)</td>
<td>{&quot;a key&quot;: &quot;1&quot;, &quot;b&quot;: &quot;t&quot;, &quot;c&quot;: null, &quot;d&quot;: &quot;12345&quot;, &quot;e&quot;: &quot;012345&quot;, &quot;f&quot;: &quot;1.234&quot;, &quot;g&quot;: &quot;2.345e+4&quot;}</td>
</tr>
<tr>
<td>hstore_to_jsonb</td>
<td>jsonb</td>
<td>get hstore as a jsonb value, converting all non-null values to JSON strings</td>
<td>hstore_to_jsonb(&quot;a key&quot;=&gt;1, &quot;b&quot;=&gt;t, &quot;c&quot;=&gt;null, &quot;d&quot;=&gt;12345, &quot;e&quot;=&gt;012345, &quot;f&quot;=&gt;1.234, &quot;g&quot;=&gt;2.345e+4&quot;)</td>
<td>{&quot;$a key&quot;: &quot;$1&quot;, &quot;$b&quot;: &quot;$t&quot;, &quot;$c&quot;: null, &quot;$d&quot;: &quot;$12345&quot;, &quot;$e&quot;: &quot;$012345&quot;, &quot;$f&quot;: &quot;$1.234&quot;, &quot;$g&quot;: &quot;$2.345e+4&quot;}</td>
</tr>
</tbody>
</table>
## Appendix F: Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
</table>
| hstore_to_json_loose(hstore)    | json        | get hstore as a json value, but attempt to distinguish numerical and Boolean values so they are unquoted in the JSON | hstore_to_json_loose("a key"=>1, "b"=>true, "c"=>null, "d"=>12345, "e": "012345", "f": 1.234, "g": 2.345e+4) | {
|                                |             |                                                                            |                                                                         |                                           |
| hstore_to_jsonb_loose(hstore)   | jsonb       | get hstore as a jsonb value, but attempt to distinguish numerical and Boolean values so they are unquoted in the JSON | hstore_to_jsonb_loose("a key"=>1, "b"=>true, "c"=>null, "d"=>12345, "e": "012345", "f": 1.234, "g": 2.345e+4) | {
|                                |             |                                                                            |                                                                         |                                           |
| slice(hstore, text[])           | hstore      | extract a subset of an hstore                                              | slice('a=>1,b=>2','a'=1,b'=2,c'=3')::hstore, ARRAY('b','c','d')=>"3" | "b"=2, "c"=3                            |
| each(hstore)                    | setof       | get hstore’s keys and values as a set                                      | select * from each('a=>1,b=>2')                                        | key | value |
|                                |             |                                                                            |                                                                         | a | 1     |
|                                |             |                                                                            |                                                                         | b | 2     |
| exist(hstore, text[])           | boolean     | does hstore contain key?                                                   | exist('a=>1','a')                                                       | true                                      |
| defined(hstore, text[])         | boolean     | does hstore contain non-NULL value for key?                                | defined('a=>NULL','a')                                                  | false                                     |
| delete(hstore, text[])          | hstore      | delete pair with matching key                                              | delete('a=>1,b=>2','a')                                                 | "a"=1, "b"=2                            |
| delete(hstore, hstore)          | hstore      | delete pairs with matching keys                                            | delete('a=>1,b=>2','a'=3,b'=4',ARRAY['a','b'])                          |                                           |
| delete(hstore, hstore[])        | hstore      | delete pairs matching those in the second argument                         | delete('a=>1,b=>2','a'=3,b'=4',ARRAY['a','b'])                          |                                           |
| populate_record(record, hstore) | record      | replace fields in record with matching values from hstore                 | see Examples section                                                   |                                           |

**Note:** The function `hstore_to_json` is used when an `hstore` value is cast to `json`. Likewise, `hstore_to_jsonb` is used when an `hstore` value is cast to `jsonb`.  

2412
Appendix F. Additional Supplied Modules

Note: The function `populate_record` is actually declared with `anyelement`, not `record`, as its first argument, but it will reject non-record types with a run-time error.

F.16.3. Indexes

`hstore` has GiST and GIN index support for the `@>`, `?`, `?&` and `?|` operators. For example:

```
CREATE INDEX hidx ON testhstore USING GIST (h);
CREATE INDEX hidx ON testhstore USING GIN (h);
```

`hstore` also supports `btree` or `hash` indexes for the `=` operator. This allows `hstore` columns to be declared `UNIQUE`, or to be used in `GROUP BY`, `ORDER BY` or `DISTINCT` expressions. The sort ordering for `hstore` values is not particularly useful, but these indexes may be useful for equivalence lookups.

Create indexes for `=` comparisons as follows:

```
CREATE INDEX hidx ON testhstore USING BTREE (h);
CREATE INDEX hidx ON testhstore USING HASH (h);
```

F.16.4. Examples

Add a key, or update an existing key with a new value:

```
UPDATE tab SET h = h || hstore('c', '3');
```

Delete a key:

```
UPDATE tab SET h = delete(h, 'k1');
```

Convert a record to an `hstore`:

```
CREATE TABLE test (col1 integer, col2 text, col3 text);
INSERT INTO test VALUES (123, 'foo', 'bar');
SELECT hstore(t) FROM test AS t;
```

```
---------------------------------------------
"col1"="123", "col2"="foo", "col3"="bar"
(1 row)
```

Convert an `hstore` to a predefined record type:

```
CREATE TABLE test (col1 integer, col2 text, col3 text);
SELECT * FROM populate_record(null::test,
    "col1"="456", "col2"="zzz");
```

```
col1 | col2 | col3
-----|------|------
2413 |      |      
```
Modify an existing record using the values from an hstore:

```sql
CREATE TABLE test (col1 integer, col2 text, col3 text);
INSERT INTO test VALUES (123, 'foo', 'bar');
SELECT (r).* FROM (SELECT t #= "col3">"baz" AS r FROM test t) s;
```

```
col1 | col2 | col3
------+------+------
123   | foo  | baz
(1 row)
```

**F.16.5. Statistics**

The hstore type, because of its intrinsic liberality, could contain a lot of different keys. Checking for valid keys is the task of the application. The following examples demonstrate several techniques for checking keys and obtaining statistics.

Simple example:

```sql
SELECT * FROM each('aaa=>bq, b=>NULL, "">1');
```

Using a table:

```sql
SELECT (each(h)).key, (each(h)).value INTO stat FROM testhstore;
```

Online statistics:

```sql
SELECT key, count(*) FROM (SELECT (each(h)).key FROM testhstore) AS stat GROUP BY key ORDER BY count DESC, key;
```

```
key   | count
------|------
line  | 883  
query | 207  
pos   | 203  
node  | 202  
space | 197  
status| 195  
pUBLIC| 194  
title | 190  
org   | 189  
```

---
F.16.6. Compatibility

As of PostgreSQL 9.0, hstore uses a different internal representation than previous versions. This presents no obstacle for dump/restore upgrades since the text representation (used in the dump) is unchanged.

In the event of a binary upgrade, upward compatibility is maintained by having the new code recognize old-format data. This will entail a slight performance penalty when processing data that has not yet been modified by the new code. It is possible to force an upgrade of all values in a table column by doing an UPDATE statement as follows:

```
UPDATE tablename SET hstorecol = hstorecol || "";
```

Another way to do it is:
```
ALTER TABLE tablename ALTER hstorecol TYPE hstore USING hstorecol || "";
```

The ALTER TABLE method requires an exclusive lock on the table, but does not result in bloating the table with old row versions.

F.16.7. Transforms

Additional extensions are available that implement transforms for the hstore type for the languages PL/Perl and PL/Python. The extensions for PL/Perl are called hstore_plperl and hstore_plperlu, for trusted and untrusted PL/Perl. If you install these transforms and specify them when creating a function, hstore values are mapped to Perl hashes. The extensions for PL/Python are called hstore_plpythonu, hstore_plpython2u, and hstore_plpython3u (see Section 44.1 for the PL/Python naming convention). If you use them, hstore values are mapped to Python dictionaries.

F.16.8. Authors

Oleg Bartunov <oleg@sai.msu.su>, Moscow, Moscow University, Russia
Teodor Sigaev <teodor@sigaev.ru>, Moscow, Delta-Soft Ltd., Russia

Additional enhancements by Andrew Gierth <andrew@taol1.riddles.org.uk>, United Kingdom

F.17. intagg

The intagg module provides an integer aggregator and an enumerator. intagg is now obsolete, because there are built-in functions that provide a superset of its capabilities. However, the module is still provided as a compatibility wrapper around the built-in functions.
F.17.1. Functions

The aggregator is an aggregate function `int_array_aggregate(integer)` that produces an integer array containing exactly the integers it is fed. This is a wrapper around `array_agg`, which does the same thing for any array type.

The enumerator is a function `int_array_enum(integer[])` that returns `setof integer`. It is essentially the reverse operation of the aggregator: given an array of integers, expand it into a set of rows. This is a wrapper around `unnest`, which does the same thing for any array type.

F.17.2. Sample Uses

Many database systems have the notion of a one to many table. Such a table usually sits between two indexed tables, for example:

```
CREATE TABLE left (id INT PRIMARY KEY, ...);
CREATE TABLE right (id INT PRIMARY KEY, ...);
CREATE TABLE one_to_many(left INT REFERENCES left, right INT REFERENCES right);
```

It is typically used like this:

```
SELECT right.* from right JOIN one_to_many ON (right.id = one_to_many.right)
WHERE one_to_many.left = item;
```

This will return all the items in the right hand table for an entry in the left hand table. This is a very common construct in SQL.

Now, this methodology can be cumbersome with a very large number of entries in the `one_to_many` table. Often, a join like this would result in an index scan and a fetch for each right hand entry in the table for a particular left hand entry. If you have a very dynamic system, there is not much you can do. However, if you have some data which is fairly static, you can create a summary table with the aggregator.

```
CREATE TABLE summary AS
    SELECT left, int_array_aggregate(right) AS right
    FROM one_to_many
GROUP BY left;
```

This will create a table with one row per left item, and an array of right items. Now this is pretty useless without some way of using the array; that’s why there is an array enumerator. You can do

```
SELECT left, int_array_enum(right) FROM summary WHERE left = item;
```

The above query using `int_array_enum` produces the same results as

```
SELECT left, right FROM one_to_many WHERE left = item;
```

The difference is that the query against the summary table has to get only one row from the table, whereas the direct query against `one_to_many` must index scan and fetch a row for each entry.

On one system, an `EXPLAIN` showed a query with a cost of 8488 was reduced to a cost of 329. The original query was a join involving the `one_to_many` table, which was replaced by:

```
SELECT right, count(right) FROM
    ( SELECT left, int_array_enum(right) AS right
    FROM summary JOIN (SELECT left FROM left_table WHERE left = item) AS lefts
    ON (summary.left = lefts.left) )
```
F.18. intarray

The intarray module provides a number of useful functions and operators for manipulating null-free arrays of integers. There is also support for indexed searches using some of the operators.

All of these operations will throw an error if a supplied array contains any NULL elements.

Many of these operations are only sensible for one-dimensional arrays. Although they will accept input arrays of more dimensions, the data is treated as though it were a linear array in storage order.

F.18.1. intarray Functions and Operators

The functions provided by the intarray module are shown in Table F-10, the operators in Table F-11.

Table F-10. intarray Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>icount(int[])</td>
<td>int</td>
<td>number of elements in array</td>
<td>icount('{1,2,3}'::int[])</td>
<td>3</td>
</tr>
<tr>
<td>sort(int[], text dir)</td>
<td>int[]</td>
<td>sort array — dir must be asc or desc</td>
<td>sort('{1,2,3}'::int[], 'desc')</td>
<td>{3,2,1}</td>
</tr>
<tr>
<td>sort(int[])</td>
<td>int[]</td>
<td>sort in ascending order</td>
<td>sort(array[11,77,44])</td>
<td>{11,44,77}</td>
</tr>
<tr>
<td>sort_asc(int[])</td>
<td>int[]</td>
<td>sort in ascending order</td>
<td></td>
<td></td>
</tr>
<tr>
<td>sort_desc(int[])</td>
<td>int[]</td>
<td>sort in descending order</td>
<td></td>
<td></td>
</tr>
<tr>
<td>uniq(int[])</td>
<td>int[]</td>
<td>remove adjacent duplicates</td>
<td>uniq(sort('{1,2,3,2,1}'::int[]))</td>
<td>{1,2,3}</td>
</tr>
<tr>
<td>idx(int[], int item)</td>
<td>int</td>
<td>index of first element matching item (0 if none)</td>
<td>idx(array[11,22,233,22,11], 22)</td>
<td>22</td>
</tr>
<tr>
<td>subarray(int[], int start, int len)</td>
<td>int[]</td>
<td>portion of array starting at position start, len elements</td>
<td>subarray('{1,2,2,3}'::int[], 2, 3)</td>
<td>{2,3}</td>
</tr>
<tr>
<td>subarray(int[], int start)</td>
<td>int[]</td>
<td>portion of array starting at position start</td>
<td>subarray('{1,2,2,3}'::int[], 2)</td>
<td>{2,3}</td>
</tr>
</tbody>
</table>
### Table F-11. `intarray` Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>int[] &amp;&amp; int[]</code></td>
<td>boolean</td>
<td>overlap — true if arrays have at least one common element</td>
</tr>
<tr>
<td><code>int[] @&gt; int[]</code></td>
<td>boolean</td>
<td>contains — true if left array contains right array</td>
</tr>
<tr>
<td><code>int[] &lt;@ int[]</code></td>
<td>boolean</td>
<td>contained — true if left array is contained in right array</td>
</tr>
<tr>
<td><code># int[]</code></td>
<td>int</td>
<td>number of elements in array</td>
</tr>
<tr>
<td><code>int[] # int</code></td>
<td>int</td>
<td>index (same as <code>idx</code> function)</td>
</tr>
<tr>
<td><code>int[] + int</code></td>
<td>int[]</td>
<td>push element onto array (add it to end of array)</td>
</tr>
<tr>
<td><code>int[] + int[]</code></td>
<td>int[]</td>
<td>array concatenation (right array added to the end of left one)</td>
</tr>
<tr>
<td><code>int[] - int</code></td>
<td>int[]</td>
<td>remove entries matching right argument from array</td>
</tr>
<tr>
<td><code>int[] - int[]</code></td>
<td>int[]</td>
<td>remove elements of right array from left</td>
</tr>
<tr>
<td>`int[]</td>
<td>int`</td>
<td>int[]</td>
</tr>
<tr>
<td>`int[]</td>
<td>int[]`</td>
<td>int[]</td>
</tr>
<tr>
<td><code>int[] &amp; int[]</code></td>
<td>int[]</td>
<td>intersection of arrays</td>
</tr>
<tr>
<td><code>int[] @@ query_int</code></td>
<td>boolean</td>
<td>true if array satisfies query (see below)</td>
</tr>
<tr>
<td><code>query_int ~~ int[]</code></td>
<td>boolean</td>
<td>true if array satisfies query (commutator of <code>@@</code>)</td>
</tr>
</tbody>
</table>

(Before PostgreSQL 8.2, the containment operators `@>` and `<@` were respectively called `@` and `~`. These names are still available, but are deprecated and will eventually be retired. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!)

The operators `&&`, `@>` and `<@` are equivalent to PostgreSQL’s built-in operators of the same names, except that they work only on integer arrays that do not contain nulls, while the built-in operators work for any array type. This restriction makes them faster than the built-in operators in many cases.

The `@@` and `~~` operators test whether an array satisfies a query, which is expressed as a value of a specialized data type `query_int`. A query consists of integer values that are checked against the elements of the array, possibly combined using the operators `&` (AND), `|` (OR), and `!` (NOT). Parentheses can be used as needed. For example, the query `1 & (2 | 3)` matches arrays that contain 1 and also contain either 2 or 3.
F.18.2. Index Support

`intarray` provides index support for the `&&`, `@>`, `<@`, and `@@` operators, as well as regular array equality.

Two GiST index operator classes are provided: `gist__int_ops` (used by default) is suitable for small- to medium-size data sets, while `gist__intbig_ops` uses a larger signature and is more suitable for indexing large data sets (i.e., columns containing a large number of distinct array values). The implementation uses an RD-tree data structure with built-in lossy compression.

There is also a non-default GIN operator class `gin__int_ops` supporting the same operators.

The choice between GiST and GIN indexing depends on the relative performance characteristics of GiST and GIN, which are discussed elsewhere.

F.18.3. Example

```
-- a message can be in one or more "sections"
CREATE TABLE message (mid INT PRIMARY KEY, sections INT[], ...);

-- create specialized index
CREATE INDEX message_rdtree_idx ON message USING GIST (sections gist__int_ops);

-- select messages in section 1 OR 2 - OVERLAP operator
SELECT message.mid FROM message WHERE message.sections && '{1,2}';

-- select messages in sections 1 AND 2 - CONTAINS operator
SELECT message.mid FROM message WHERE message.sections @> '{1,2}';

-- the same, using QUERY operator
SELECT message.mid FROM message WHERE message.sections @@ '1&2':query_int;
```

F.18.4. Benchmark

The source directory `contrib/intarray/bench` contains a benchmark test suite, which can be run against an installed PostgreSQL server. (It also requires DBD::Pg to be installed.) To run:

```
cd .../contrib/intarray/bench
createdb TEST
psql -c "CREATE EXTENSION intarray" TEST
./create_test.pl | psql TEST
./bench.pl
```

The `bench.pl` script has numerous options, which are displayed when it is run without any arguments.

F.18.5. Authors

All work was done by Teodor Sigaev (<teodor@sigaev.ru>) and Oleg Bartunov (<oleg@sai.msu.su>). See http://www.sai.msu.su/~megera/postgres/gist/ for additional information. Andrey Oktyabrski did a great work on adding new functions and operations.
F.19. isn

The isn module provides data types for the following international product numbering standards: EAN13, UPC, ISBN (books), ISMN (music), and ISSN (serials). Numbers are validated on input according to a hard-coded list of prefixes; this list of prefixes is also used to hyphenate numbers on output. Since new prefixes are assigned from time to time, the list of prefixes may be out of date. It is hoped that a future version of this module will obtained the prefix list from one or more tables that can be easily updated by users as needed; however, at present, the list can only be updated by modifying the source code and recompiling. Alternatively, prefix validation and hyphenation support may be dropped from a future version of this module.

F.19.1. Data Types

Table F-12 shows the data types provided by the isn module.

<table>
<thead>
<tr>
<th>Data Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>EAN13</td>
<td>European Article Numbers, always displayed in the EAN13 display format</td>
</tr>
<tr>
<td>ISBN13</td>
<td>International Standard Book Numbers to be displayed in the new EAN13 display format</td>
</tr>
<tr>
<td>ISMN13</td>
<td>International Standard Music Numbers to be displayed in the new EAN13 display format</td>
</tr>
<tr>
<td>ISSN13</td>
<td>International Standard Serial Numbers to be displayed in the new EAN13 display format</td>
</tr>
<tr>
<td>ISBN</td>
<td>International Standard Book Numbers to be displayed in the old short display format</td>
</tr>
<tr>
<td>ISMN</td>
<td>International Standard Music Numbers to be displayed in the old short display format</td>
</tr>
<tr>
<td>ISSN</td>
<td>International Standard Serial Numbers to be displayed in the old short display format</td>
</tr>
<tr>
<td>UPC</td>
<td>Universal Product Codes</td>
</tr>
</tbody>
</table>

Some notes:

1. ISBN13, ISMN13, ISSN13 numbers are all EAN13 numbers.
2. EAN13 numbers aren’t always ISBN13, ISMN13 or ISSN13 (some are).
3. Some ISBN13 numbers can be displayed as ISBN.
4. Some ISMN13 numbers can be displayed as ISMN.
5. Some ISSN13 numbers can be displayed as ISSN.
6. UPC numbers are a subset of the EAN13 numbers (they are basically EAN13 without the first 0 digit).
7. All UPC, ISBN, ISMN and ISSN numbers can be represented as EAN13 numbers.

Internally, all these types use the same representation (a 64-bit integer), and all are interchangeable. Multiple types are provided to control display formatting and to permit tighter validity checking of input that is supposed to denote one particular type of number.
The ISBN, ISMN, and ISSN types will display the short version of the number (ISxN 10) whenever it’s possible, and will show ISxN 13 format for numbers that do not fit in the short version. The EAN13, ISBN13, ISMN13 and ISSN13 types will always display the long version of the ISxN (EAN13).

F.19.2. Casts

The isn module provides the following pairs of type casts:

- ISBN13 <= > EAN13
- ISMN13 <= > EAN13
- ISSN13 <= > EAN13
- ISBN <= > EAN13
- ISMN <= > EAN13
- ISSN <= > EAN13
- UPC <= > EAN13
- ISMN <= > ISMN13
- ISSN <= > ISSN13

When casting from EAN13 to another type, there is a run-time check that the value is within the domain of the other type, and an error is thrown if not. The other casts are simply relabelings that will always succeed.

F.19.3. Functions and Operators

The isn module provides the standard comparison operators, plus B-tree and hash indexing support for all these data types. In addition there are several specialized functions; shown in Table F-13. In this table, isn means any one of the module’s data types.

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>isn_weak(boolean)</td>
<td>boolean</td>
<td>Sets the weak input mode (returns new setting)</td>
</tr>
<tr>
<td>isn_weak()</td>
<td>boolean</td>
<td>Gets the current status of the weak mode</td>
</tr>
<tr>
<td>make_valid(isn)</td>
<td>isn</td>
<td>Validates an invalid number (clears the invalid flag)</td>
</tr>
<tr>
<td>is_valid(isn)</td>
<td>boolean</td>
<td>Checks for the presence of the invalid flag</td>
</tr>
</tbody>
</table>

Weak mode is used to be able to insert invalid data into a table. Invalid means the check digit is wrong, not that there are missing numbers.

Why would you want to use the weak mode? Well, it could be that you have a huge collection of ISBN numbers, and that there are so many of them that for weird reasons some have the wrong check digit (perhaps the numbers were scanned from a printed list and the OCR got the numbers wrong,
perhaps the numbers were manually captured... who knows). Anyway, the point is you might want to clean the mess up, but you still want to be able to have all the numbers in your database and maybe use an external tool to locate the invalid numbers in the database so you can verify the information and validate it more easily; so for example you’d want to select all the invalid numbers in the table.

When you insert invalid numbers in a table using the weak mode, the number will be inserted with the corrected check digit, but it will be displayed with an exclamation mark (!) at the end, for example 0-11-000322-5!. This invalid marker can be checked with the is_valid function and cleared with the make_valid function.

You can also force the insertion of invalid numbers even when not in the weak mode, by appending the ! character at the end of the number.

Another special feature is that during input, you can write ? in place of the check digit, and the correct check digit will be inserted automatically.

### F.19.4. Examples

--Using the types directly:
```
SELECT isbn('978-0-393-04002-9');
SELECT isbn13('0901690546');
SELECT issn('1436-4522');
```

--Casting types:
```
-- note that you can only cast from ean13 to another type when the
-- number would be valid in the realm of the target type;
-- thus, the following will NOT work: select isbn(ean13('0220356483481'));
-- but these will:
SELECT upc(ean13('0220356483481'));
SELECT ean13(upc('220356483481'));
```

--Create a table with a single column to hold ISBN numbers:
```
CREATE TABLE test (id isbn);
INSERT INTO test VALUES('9780393040029');
--Automatically calculate check digits (observe the '?'):
INSERT INTO test VALUES('220500896?');
INSERT INTO test VALUES('978055215372?');
SELECT issn('3251231?');
SELECT ismn('979047213542?');
```

--Using the weak mode:
```
SELECT isn_weak(true);
INSERT INTO test VALUES('978-0-11-000533-4');
INSERT INTO test VALUES('9780141219307');
INSERT INTO test VALUES('2-205-00876-X');
SELECT isn_weak(false);
SELECT id FROM test WHERE NOT is_valid(id);
UPDATE test SET id = make_valid(id) WHERE id = '2-205-00876-X!';
SELECT * FROM test;
SELECT isbn13(id) FROM test;
```
F.19.5. Bibliography

The information to implement this module was collected from several sites, including:

- http://www.isbn-international.org/
- http://www.issn.org/
- http://www.ismn-international.org/
- http://www.wikipedia.org/

The prefixes used for hyphenation were also compiled from:

- http://www.gs1.org/productssolutions/idkeys/support/prefix_list.html
- https://www.isbn-international.org/content/isbn-users-manual
- http://www.ismn-international.org/ranges.html

Care was taken during the creation of the algorithms and they were meticulously verified against the suggested algorithms in the official ISBN, ISMN, ISSN User Manuals.

F.19.6. Author

Germán Méndez Bravo (Kronuz), 2004 - 2006

This module was inspired by Garrett A. Wollman’s isbn_issn code.

F.20. lo

The lo module provides support for managing Large Objects (also called LOs or BLOBs). This includes a data type lo and a trigger lo_manage.

F.20.1. Rationale

One of the problems with the JDBC driver (and this affects the ODBC driver also), is that the specification assumes that references to BLOBs (Binary Large OBjects) are stored within a table, and if that entry is changed, the associated BLOB is deleted from the database.

As PostgreSQL stands, this doesn’t occur. Large objects are treated as objects in their own right; a table entry can reference a large object by OID, but there can be multiple table entries referencing the same large object OID, so the system doesn’t delete the large object just because you change or remove one such entry.

Now this is fine for PostgreSQL-specific applications, but standard code using JDBC or ODBC won’t delete the objects, resulting in orphan objects — objects that are not referenced by anything, and simply occupy disk space.

The lo module allows fixing this by attaching a trigger to tables that contain LO reference columns. The trigger essentially just does a lo_unlink whenever you delete or modify a value referencing a
large object. When you use this trigger, you are assuming that there is only one database reference to any large object that is referenced in a trigger-controlled column!

The module also provides a data type `lo`, which is really just a domain of the `oid` type. This is useful for differentiating database columns that hold large object references from those that are OIDs of other things. You don’t have to use the `lo` type to use the trigger, but it may be convenient to use it to keep track of which columns in your database represent large objects that you are managing with the trigger. It is also rumored that the ODBC driver gets confused if you don’t use `lo` for BLOB columns.

**F.20.2. How to Use It**

Here’s a simple example of usage:

```sql
CREATE TABLE image (title TEXT, raster lo);

CREATE TRIGGER t_raster BEFORE UPDATE OR DELETE ON image
    FOR EACH ROW EXECUTE PROCEDURE lo_manage(raster);
```

For each column that will contain unique references to large objects, create a `BEFORE UPDATE OR DELETE` trigger, and give the column name as the sole trigger argument. You can also restrict the trigger to only execute on updates to the column by using `BEFORE UPDATE OF column_name`. If you need multiple `lo` columns in the same table, create a separate trigger for each one, remembering to give a different name to each trigger on the same table.

**F.20.3. Limitations**

- Dropping a table will still orphan any objects it contains, as the trigger is not executed. You can avoid this by preceding the `DROP TABLE` with `DELETE FROM table`. `TRUNCATE` has the same hazard.

- If you already have, or suspect you have, orphaned large objects, see the `vacuumlo` module to help you clean them up. It’s a good idea to run `vacuumlo` occasionally as a back-stop to the `lo_manage` trigger.

- Some frontends may create their own tables, and will not create the associated trigger(s). Also, users may not remember (or know) to create the triggers.

**F.20.4. Author**

Peter Mount <peter@retep.org.uk>

**F.21. ltree**

This module implements a data type `ltree` for representing labels of data stored in a hierarchical tree-like structure. Extensive facilities for searching through label trees are provided.
F.21.1. Definitions

A **label** is a sequence of alphanumeric characters and underscores (for example, in C locale the characters A-Za-z0-9_ are allowed). Labels must be less than 256 bytes long.

**Examples:** 42, Personal_Services

A **label path** is a sequence of zero or more labels separated by dots, for example L1.L2.L3, representing a path from the root of a hierarchical tree to a particular node. The length of a label path must be less than 65kB, but keeping it under 2kB is preferable.

**Example:** Top.Countries.Europe.Russia

The **ltree** module provides several data types:

- **ltree** stores a label path.
- **lquery** represents a regular-expression-like pattern for matching ltree values. A simple word matches that label within a path. A star symbol (*) matches zero or more labels. For example:

```
foo          Match the exact label path foo
*.foo.*      Match any label path containing the label foo
*.foo        Match any label path whose last label is foo
```

Star symbols can also be quantified to restrict how many labels they can match:

- `*(n)` Match exactly n labels
- `*(n),` Match at least n labels
- `*(n,m)` Match at least n but not more than m labels
- `*(,m)` Match at most m labels -- same as `*(0,m)`

There are several modifiers that can be put at the end of a non-star label in lquery to make it match more than just the exact match:

- `@` Match case-insensitively, for example a@ matches A
- `*` Match any label with this prefix, for example foo* matches foobar
- `%` Match initial underscore-separated words

The behavior of `%` is a bit complicated. It tries to match words rather than the entire label. For example foo_bar% matches foo_bar baz but not foo_bar baz. If combined with *, prefix matching applies to each word separately, for example foo_bar%* matches foo1_bar2_baz but not foo1_br2_baz.

Also, you can write several possibly-modified labels separated with | (OR) to match any of those labels, and you can put ! (NOT) at the start to match any label that doesn’t match any of the alternatives.

Here’s an annotated example of lquery:
```
Top.*(0,2).sport@!.football|tennis.Russ*|Spain
```

This query will match any label path that:

a. begins with the label Top
b. and next has zero to two labels before
c. a label beginning with the case-insensitive prefix sport
d. then a label not matching football nor tennis
e. and then ends with a label beginning with Russ or exactly matching Spain.
Appendix F. Additional Supplied Modules

• `ltxtquery` represents a full-text-search-like pattern for matching `ltree` values. An `ltxtquery` value contains words, possibly with the modifiers `@`, `*`, `%` at the end; the modifiers have the same meanings as in `lquery`. Words can be combined with `&` (AND), `|` (OR), `!` (NOT), and parentheses. The key difference from `lquery` is that `ltxtquery` matches words without regard to their position in the label path.

Here’s an example `ltxtquery`:

`Europe & Russia* & !Transportation`

This will match paths that contain the label `Europe` and any label beginning with `Russia` (case-insensitive), but not paths containing the label `Transportation`. The location of these words within the path is not important. Also, when `%` is used, the word can be matched to any underscore-separated word within a label, regardless of position.

Note: `ltxtquery` allows whitespace between symbols, but `ltree` and `lquery` do not.

F.21.2. Operators and Functions

Type `ltree` has the usual comparison operators `=`, `<`, `<=`, `>=`, `>`. Comparison sorts in the order of a tree traversal, with the children of a node sorted by label text. In addition, the specialized operators shown in Table F-14 are available.

Table F-14. `ltree` Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>ltree &gt;= ltree</code></td>
<td>boolean</td>
<td>is left argument an ancestor of right (or equal)?</td>
</tr>
<tr>
<td><code>ltree &lt;= ltree</code></td>
<td>boolean</td>
<td>is left argument a descendant of right (or equal)?</td>
</tr>
<tr>
<td><code>ltree ~ lquery</code></td>
<td>boolean</td>
<td>does <code>ltree</code> match <code>lquery</code>?</td>
</tr>
<tr>
<td><code>lquery ~ ltree</code></td>
<td>boolean</td>
<td>does <code>ltree</code> match <code>lquery</code>?</td>
</tr>
<tr>
<td><code>ltree ? lquery[]</code></td>
<td>boolean</td>
<td>does <code>ltree</code> match any <code>lquery</code> in array?</td>
</tr>
<tr>
<td><code>lquery[] ? ltree</code></td>
<td>boolean</td>
<td>does <code>ltree</code> match any <code>lquery</code> in array?</td>
</tr>
<tr>
<td>boolean</td>
<td>does <code>ltree</code> match <code>lquery</code>?</td>
<td></td>
</tr>
<tr>
<td><code>ltree @ lquery</code></td>
<td>boolean</td>
<td>does <code>ltree</code> match <code>lquery</code>?</td>
</tr>
<tr>
<td><code>ltree @ ltxtquery</code></td>
<td>boolean</td>
<td>does <code>ltree</code> match <code>ltxtquery</code>?</td>
</tr>
<tr>
<td><code>ltxtquery @ ltree</code></td>
<td>boolean</td>
<td>does <code>ltxtquery</code> match <code>ltree</code>?</td>
</tr>
<tr>
<td>`ltree</td>
<td></td>
<td>ltree`</td>
</tr>
<tr>
<td>`ltree</td>
<td></td>
<td>text`</td>
</tr>
<tr>
<td>`text</td>
<td></td>
<td>ltree`</td>
</tr>
<tr>
<td><code>ltree[] @&gt; ltree</code></td>
<td>boolean</td>
<td>does array contain an ancestor of <code>ltree</code>?</td>
</tr>
<tr>
<td><code>ltree &lt;= ltree[]</code></td>
<td>boolean</td>
<td>does array contain an ancestor of <code>ltree</code>?</td>
</tr>
</tbody>
</table>
### Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ltree[] &lt; ltree</td>
<td>boolean</td>
<td>does array contain a descendant of ltree?</td>
</tr>
<tr>
<td>ltree @ ltree[]</td>
<td>boolean</td>
<td>does array contain a descendant of ltree?</td>
</tr>
<tr>
<td>ltree[] ~ lquery</td>
<td>boolean</td>
<td>does array contain any path matching lquery?</td>
</tr>
<tr>
<td>lquery ~ ltree[]</td>
<td>boolean</td>
<td>does array contain any path matching lquery?</td>
</tr>
<tr>
<td>ltree[] ? lquery[]</td>
<td>boolean</td>
<td>does ltree array contain any path matching lquery?</td>
</tr>
<tr>
<td>lquery[] ? ltree[]</td>
<td>boolean</td>
<td>does ltree array contain any path matching lquery?</td>
</tr>
<tr>
<td>ltree[] @ ltquery</td>
<td>boolean</td>
<td>does array contain any path matching ltquery?</td>
</tr>
<tr>
<td>ltquery @ ltree[]</td>
<td>boolean</td>
<td>does array contain any path matching ltquery?</td>
</tr>
<tr>
<td>ltree[] ? @ ltree</td>
<td>ltree</td>
<td>first array entry that is an ancestor of ltree; NULL if none</td>
</tr>
<tr>
<td>ltree[] ? &lt; @ ltree</td>
<td>ltree</td>
<td>first array entry that is a descendant of ltree; NULL if none</td>
</tr>
<tr>
<td>ltree[] ? ~ lquery</td>
<td>ltree</td>
<td>first array entry that matches lquery; NULL if none</td>
</tr>
<tr>
<td>ltree[] ? @ ltquery</td>
<td>ltree</td>
<td>first array entry that matches ltquery; NULL if none</td>
</tr>
</tbody>
</table>

The operators <, @, and ~ have analogues ^<, ^@, and ^~, which are the same except they do not use indexes. These are useful only for testing purposes.

The available functions are shown in Table F-15.

### Table F-15. ltree Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
</table>
| subltree(ltree, int start, int end) | ltree      | subpath of ltree from position start to position end-1 (counting from 0) | subltree('Top.Child1.Child2',1,2) | }
<table>
<thead>
<tr>
<th>Function</th>
<th>Return Type</th>
<th>Description</th>
<th>Example</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>subpath(ltree, int offset, int len)</td>
<td>ltree</td>
<td>subpath of ltree starting at position offset, length len. If offset is negative, subpath starts that far from the end of the path. If len is negative, leaves that many labels off the end of the path.</td>
<td>subpath('Top.Ch1.Ch2.Ch3',0,2)</td>
<td>subpath('Top.Ch1.Ch2.Ch3',0,2)</td>
</tr>
<tr>
<td>subpath(ltree, int offset)</td>
<td>ltree</td>
<td>subpath of ltree starting at position offset, extending to end of path. If offset is negative, subpath starts that far from the end of the path.</td>
<td>subpath('Top.Ch1.Ch2.Ch3',1)</td>
<td>subpath('Top.Ch1.Ch2.Ch3',1)</td>
</tr>
<tr>
<td>nlevel(ltree)</td>
<td>integer</td>
<td>number of labels in path</td>
<td>nlevel('Top.Ch1.Ch2.Ch3')</td>
<td>nlevel('Top.Ch1.Ch2.Ch3')</td>
</tr>
<tr>
<td>index(ltree a, ltree b)</td>
<td>integer</td>
<td>position of first occurrence of b in a: -1 if not found</td>
<td>index('0.1.2.3.4.5.6.7.8','5.6')</td>
<td>index('0.1.2.3.4.5.6.7.8','5.6')</td>
</tr>
<tr>
<td>index(ltree a, ltree b, int offset)</td>
<td>integer</td>
<td>position of first occurrence of b in a, searching starting at offset; negative offset means start -offset labels from the end of the path</td>
<td>index('0.1.2.3.4.5.6.7.8','5.6',-4)</td>
<td>index('0.1.2.3.4.5.6.7.8','5.6',-4)</td>
</tr>
<tr>
<td>text2ltree(text)</td>
<td>ltree</td>
<td>cast text to ltree</td>
<td>text2ltree('Top.Ch1.Ch2.Ch3')</td>
<td>text2ltree('Top.Ch1.Ch2.Ch3')</td>
</tr>
<tr>
<td>ltree2text(ltree)</td>
<td>text</td>
<td>cast ltree to text</td>
<td>ltree2text('Top.Ch1.Ch2.Ch3')</td>
<td>ltree2text('Top.Ch1.Ch2.Ch3')</td>
</tr>
<tr>
<td>lca(ltree, ltree, ...)</td>
<td>ltree</td>
<td>longest common ancestor of paths (up to 8 arguments supported)</td>
<td>lca('1.2.3','1.2.4.5.6')</td>
<td>lca('1.2.3','1.2.4.5.6')</td>
</tr>
<tr>
<td>lca(ltree[])</td>
<td>ltree</td>
<td>longest common ancestor of paths in array</td>
<td>lca(array['1.2,3:ltree', '1.2.3.4'])</td>
<td>lca(array['1.2,3:ltree', '1.2.3.4'])</td>
</tr>
</tbody>
</table>
F.21.3. Indexes

`ltree` supports several types of indexes that can speed up the indicated operators:

- **B-tree index over `ltree`: `<, <=, =, >, >=`
- **GiST index over `ltree`: `<, <=, =, >, >=`

Example of creating such an index:

```
CREATE INDEX path_gist_idx ON test USING GIST (path);
```

- **GiST index over `ltree[]`: `ltree[] <, <=, =, >, @, >`, `ltree[] >`, `ltree[] <, @, ~, ?`**

Example of creating such an index:

```
CREATE INDEX path_gist_idx ON test USING GIST (array_path);
```

Note: This index type is lossy.

F.21.4. Example

This example uses the following data (also available in file `contrib/ltree/ltreetest.sql` in the source distribution):

```
CREATE TABLE test (path ltree);
INSERT INTO test VALUES ('Top');
INSERT INTO test VALUES ('Top.Science');
INSERT INTO test VALUES ('Top.Science.Astronomy');
INSERT INTO test VALUES ('Top.Science.Astronomy.Astrophysics');
INSERT INTO test VALUES ('Top.Science.Astronomy.Cosmology');
INSERT INTO test VALUES ('Top.Hobbies');
INSERT INTO test VALUES ('Top.Hobbies.Amateurs_Astronomy');
INSERT INTO test VALUES ('Top.Collections');
INSERT INTO test VALUES ('Top.Collections.Pictures');
INSERT INTO test VALUES ('Top.Collections.Pictures.Astronomy');
INSERT INTO test VALUES ('Top.Collections.Pictures.Astronomy.Astronauts');
CREATE INDEX path_gist_idx ON test USING GIST (path);
CREATE INDEX path_idx ON test USING BTREE (path);
```

Now, we have a table `test` populated with data describing the hierarchy shown below:

```
<table>
<thead>
<tr>
<th>Top</th>
<th>Science</th>
<th>Hobbies</th>
<th>Collections</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Astronomy</td>
<td>Amateurs_Astronomy</td>
<td>Pictures</td>
</tr>
<tr>
<td></td>
<td>Astrophysics</td>
<td>Cosmology</td>
<td>Astronomy</td>
</tr>
<tr>
<td></td>
<td>Galaxy</td>
<td>Stars</td>
<td>Astronauts</td>
</tr>
</tbody>
</table>
```

We can do inheritance:

```
ltree=> SELECT path FROM test WHERE path @ 'Top.Science';
```

```
path
2429
```

2429
Here are some examples of path matching:

```
lrtree=\$\>\$ SELECT path FROM test WHERE path ~ '.*.Astronomy.*';
       path
  Top.Science.Astronomy
  Top.Science.Astronomy.Astrophysics
  Top.Collections.Pictures.Astronomy
  (7 rows)
```

```
lrtree=\$\>\$ SELECT path FROM test WHERE path ~ '*.!pictures@.*.Astronomy.*';
       path
  Top.Science.Astronomy
  Top.Science.Astronomy.Astrophysics
  (3 rows)
```

Here are some examples of full text search:

```
lrtree=\$\>\$ SELECT path FROM test WHERE path @ 'Astro*% & !pictures@';
       path
  Top.Science.Astronomy
  Top.Science.Astronomy.Astrophysics
  Top.Hobbies.Amateurs_Astronomy
  (4 rows)
```

```
lrtree=\$\>\$ SELECT path FROM test WHERE path @ 'Astro* & !pictures@';
       path
  Top.Science.Astronomy
  Top.Science.Astronomy.Astrophysics
  (3 rows)
```

Path construction using functions:

```
lrtree=\$\>\$ SELECT subpath(path,0,2)||'Space'||subpath(path,2) FROM test WHERE path <\0 ?column?
                    path
  Top.Science.Space.Astronomy
```

2430
We could simplify this by creating a SQL function that inserts a label at a specified position in a path:

```sql
CREATE FUNCTION ins_label(ltree, int, text) RETURNS ltree
AS 'select subpath($1,0,$2) || $3 || subpath($1,$2);'
LANGUAGE SQL IMMUTABLE;
```

```sql
ltreetest=> SELECT ins_label(path,2,'Space') FROM test WHERE path @@ 'Top.Science.Astronomy';
```

<table>
<thead>
<tr>
<th>ins_label</th>
</tr>
</thead>
<tbody>
<tr>
<td>Top.Science.Space.Astronomy</td>
</tr>
</tbody>
</table>

F.21.5. Transforms

Additional extensions are available that implement transforms for the `ltree` type for PL/Python. The extensions are called `ltree_plpythonu`, `ltree_plpython2u`, and `ltree_plpython3u` (see Section 44.1 for the PL/Python naming convention). If you install these transforms and specify them when creating a function, `ltree` values are mapped to Python lists. (The reverse is currently not supported, however.)

F.21.6. Authors

All work was done by Teodor Sigaev (<teodor@stack.net>) and Oleg Bartunov (<oleg@sai.msu.su>). See [http://www.sai.msu.su/~megera/postgres/gist/](http://www.sai.msu.su/~megera/postgres/gist/) for additional information. Authors would like to thank Eugeny Rodichev for helpful discussions. Comments and bug reports are welcome.

F.22. paigeinspect

The `paigeinspect` module provides functions that allow you to inspect the contents of database pages at a low level, which is useful for debugging purposes. All of these functions may be used only by superusers.

F.22.1. Functions

```sql
get_raw_page(relname text, fork text, blkno int) returns bytea
```

`get_raw_page` reads the specified block of the named relation and returns a copy as a `bytea` value. This allows a single time-consistent copy of the block to be obtained. `fork` should be
Appendix F. Additional Supplied Modules

‘main’ for the main data fork, ‘fsm’ for the free space map, ‘vm’ for the visibility map, or ‘init’ for the initialization fork.

get_raw_page(relname text, blkno int) returns bytea

A shorthand version of get_raw_page, for reading from the main fork. Equivalent to get_raw_page(relname, ‘main’, blkno)

page_header(page bytea) returns record

page_header shows fields that are common to all PostgreSQL heap and index pages.

A page image obtained with get_raw_page should be passed as argument. For example:

test=# SELECT * FROM page_header(get_raw_page('pg_class', 0));

lsn | checksum | flags | lower | upper | special | pagesize | version | prune_xid
---+----------+--------+-------+-------+---------+----------+---------+-----------
0/24A1B50 | 1 | 1 | 232 | 368 | 8192 | 8192 | 4 | 0

The returned columns correspond to the fields in the PageHeaderData struct. See src/include/storage/bufpage.h for details.

heap_page_items(page bytea) returns setof record

heap_page_items shows all line pointers on a heap page. For those line pointers that are in use, tuple headers as well as tuple raw data are also shown. All tuples are shown, whether or not the tuples were visible to an MVCC snapshot at the time the raw page was copied.

A heap page image obtained with get_raw_page should be passed as argument. For example:

test=# SELECT * FROM heap_page_items(get_raw_page('pg_class', 0));

See src/include/storage/itemid.h and src/include/access/htup_details.h for explanations of the fields returned.

tuple_data_split(rel_oid oid, t_data bytea, t_infomask integer, t_infomask2 integer, t_bits text [, do_detoast bool]) returns bytea[]

tuple_data_split splits tuple data into attributes in the same way as backend internals.

This function should be called with the same arguments as the return attributes of heap_page_items.

If do_detoast is true, attribute that will be detoasted as needed. Default value is false.

heap_page_item_attrs(page bytea, rel_oid regclass [, do_detoast bool])
returns setof record

heap_page_item_attrs is equivalent to heap_page_items except that it returns tuple raw data as an array of attributes that can optionally be detoasted by do_detoast which is false by default.

A heap page image obtained with get_raw_page should be passed as argument. For example:

bt_metap(relname text) returns record

bt_metap returns information about a B-tree index’s metapage. For example:

test=# SELECT * FROM bt_metap('pg_cast_oid_index');

-[ RECORD 1 ]-----
magic | 340322
version | 2
root | 1
level | 0
fastroot | 1

2432
Appendix F. Additional Supplied Modules

fastlevel | 0

bt_page_stats(relname text, blkno int) returns record

bt_page_stats returns summary information about single pages of B-tree indexes. For example:

test=# SELECT * FROM bt_page_stats('pg_cast_oid_index', 1);

- [ RECORD 1 ]-+-----
| blkno | 1 |
| type | l |
| live_items | 256 |
| dead_items | 0 |
| avg_item_size | 12 |
| page_size | 8192 |
| free_size | 4056 |
| btpo_prev | 0 |
| btpo_next | 0 |
| btpo | 0 |
| btpo_flags | 3 |

test_page_items(relname text, blkno int) returns setof record

bt_page_items returns detailed information about all of the items on a B-tree index page. For example:

test=# SELECT * FROM bt_page_items('pg_cast_oid_index', 1);

| itemoffset | ctid | itemlen | nulls | vars | data |
|------------+---------+---------+-------+------+--------|
| 1 | (0,1) | 12 | f | f | 23 27 00 00 |
| 2 | (0,2) | 12 | f | f | 24 27 00 00 |
| 3 | (0,3) | 12 | f | f | 25 27 00 00 |
| 4 | (0,4) | 12 | f | f | 26 27 00 00 |
| 5 | (0,5) | 12 | f | f | 27 27 00 00 |
| 6 | (0,6) | 12 | f | f | 28 27 00 00 |
| 7 | (0,7) | 12 | f | f | 29 27 00 00 |
| 8 | (0,8) | 12 | f | f | 2a 27 00 00 |

In a B-tree leaf page, ctid points to a heap tuple. In an internal page, the block number part of ctid points to another page in the index itself, while the offset part (the second number) is ignored and is usually 1.

Note that the first item on any non-rightmost page (any page with a non-zero value in the btpo_next field) is the page’s “high key”, meaning its data serves as an upper bound on all items appearing on the page, while its ctid field is meaningless. Also, on non-leaf pages, the first real data item (the first item that is not a high key) is a “minus infinity” item, with no actual value in its data field. Such an item does have a valid downlink in its ctid field, however.

brin_page_type(page bytea) returns text

brin_page_type returns the page type of the given BRIN index page, or throws an error if the page is not a valid BRIN page. For example:

test=# SELECT brin_page_type(get_raw_page('brinidx', 0));

<table>
<thead>
<tr>
<th>brin_page_type</th>
</tr>
</thead>
<tbody>
<tr>
<td>meta</td>
</tr>
</tbody>
</table>

2433
Appendix F. Additional Supplied Modules

brin_metapage_info(page bytea) returns record

brin_metapage_info returns assorted information about a BRIN index metapage. For example:

test=# SELECT * FROM brin_metapage_info(get_raw_page('brinidx', 0));
    magic | version | pagesperrange | lastrevmappage
----------+---------+---------------+----------------
   0xA8109CFA |   1 |               4 |          2

brin_revmap_data(page bytea) returns setof tid

brin_revmap_data returns the list of tuple identifiers in a BRIN index range map page. For example:

test=# SELECT * FROM brin_revmap_data(get_raw_page('brinidx', 2)) limit 5;
    pages
----------
   (6,137)
   (6,138)
   (6,139)
   (6,140)
   (6,141)

brin_page_items(page bytea, index oid) returns setof record

brin_page_items returns the data stored in the BRIN data page. For example:

test=# SELECT * FROM brin_page_items(get_raw_page('brinidx', 5),
                      'brinidx')
    ORDER BY blknum, attnum LIMIT 6;
    itemoffset | blknum | attnum | allnulls | hasnulls | placeholder | value
----------+--------+--------+----------+----------+-------------+--------------
    137 |   0 |       1 |        t |         f |            f | {1 .. 88}
    138 |   4 |       1 |        t |         f |            f | {89 .. 176}
    139 |   8 |       1 |        t |         f |            f |  
    139 |   8 |       2 |        f |         f |            f | {177 .. 264}

The returned columns correspond to the fields in the BrinMemTuple and BrinValues structs.
See src/include/access/brin_tuple.h for details.

gin_metapage_info(page bytea) returns record

gin_metapage_info returns information about a GIN index metapage. For example:

test=# SELECT * FROM gin_metapage_info(get_raw_page('gin_index', 0));

-{ RECORD 1 }----------------------------------------
   pending_head | 4294967295
   pending_tail | 4294967295
tail_free_size | 0
n_pending_pages | 0
n_pending_tuples | 0
n_total_pages | 7
n_entry_pages | 6
n_data_pages | 0
n_entries | 693
   version | 2

2434
gin_page_opaque_info(page bytea) returns record

gin_page_opaque_info returns information about a GIN index opaque area, like the page type. For example:

test=# SELECT * FROM gin_page_opaque_info(get_raw_page('gin_index', 2));

rightlink | maxoff | flags
-----------+--------+------------------------
      5 |      0 | {data,leaf,compressed}
(1 row)

gin_leafpage_items(page bytea) returns setof record

gin_leafpage_items returns information about the data stored in a GIN leaf page. For example:

test=# SELECT first_tid, nbytes, tids[0:5] as some_tids
FROM gin_leafpage_items(get_raw_page('gin_test_idx', 2));

first_tid | nbytes | some_tids
-----------+--------+----------------------------------------------------------
     (8,41) |  244 | {"(8,41)"", "(8,43)"", "(8,44)"", "(8,45)"", "(8,46)"}
  (10,45) |  248 | {"(10,45)"", "(10,46)"", "(10,47)"", "(10,48)"", "(10,49)"}
   (12,52) |  248 | {"(12,52)"", "(12,53)"", "(12,54)"", "(12,55)"", "(12,56)"}
   (14,59) |  320 | {"(14,59)"", "(14,60)"", "(14,61)"", "(14,62)"", "(14,63)"}
  (167,16) |  376 | {"(167,16)"", "(167,17)"", "(167,18)"", "(167,19)"", "(167,20)"}
 (170,30) |  376 | {"(170,30)"", "(170,31)"", "(170,32)"", "(170,33)"", "(170,34)"}
 (173,44) |  197 | {"(173,44)"", "(173,45)"", "(173,46)"", "(173,47)"", "(173,48)"}
(7 rows)

fsm_page_contents(page bytea) returns text

fsm_page_contents shows the internal node structure of a FSM page. The output is a multiline string, with one line per node in the binary tree within the page. Only those nodes that are not zero are printed. The so-called "next" pointer, which points to the next slot to be returned from the page, is also printed.

See src/backend/storage/freespace/README for more information on the structure of an FSM page.

F.23. passwordcheck

The passwordcheck module checks users' passwords whenever they are set with CREATE ROLE or ALTER ROLE. If a password is considered too weak, it will be rejected and the command will terminate with an error.

To enable this module, add '$libdir/passwordcheck' to shared_preload_libraries in postgresql.conf, then restart the server.

You can adapt this module to your needs by changing the source code. For example, you can use CrackLib to check passwords — this only requires uncommenting two lines in the Makefile and rebuilding the module. (We cannot include CrackLib by default for license reasons.) Without CrackLib, the module enforces a few simple rules for password strength, which you can modify or extend as you see fit.

Caution
To prevent unencrypted passwords from being sent across the network, written to the server log or otherwise stolen by a database administrator, PostgreSQL allows the user to supply pre-encrypted passwords. Many client programs make use of this functionality and encrypt the password before sending it to the server.

This limits the usefulness of the `passwordcheck` module, because in that case it can only try to guess the password. For this reason, `passwordcheck` is not recommended if your security requirements are high. It is more secure to use an external authentication method such as GSSAPI (see Chapter 20) than to rely on passwords within the database.

Alternatively, you could modify `passwordcheck` to reject pre-encrypted passwords, but forcing users to set their passwords in clear text carries its own security risks.

F.24. pg_buffercache

The `pg_buffercache` module provides a means for examining what’s happening in the shared buffer cache in real time.

The module provides a C function `pg_buffercache_pages` that returns a set of records, plus a view `pg_buffercache` that wraps the function for convenient use.

By default public access is revoked from both of these, just in case there are security issues lurking.

F.24.1. The pg_buffercache View

The definitions of the columns exposed by the view are shown in Table F-16.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bufferid</td>
<td>integer</td>
<td></td>
<td>ID, in the range 1..<code>shared_buffers</code></td>
</tr>
<tr>
<td>relfilenode</td>
<td>oid</td>
<td><code>pg_class.relfilenode</code></td>
<td>Filenode number of the relation</td>
</tr>
<tr>
<td>reltablespace</td>
<td>oid</td>
<td><code>pg_tablespace.oid</code></td>
<td>Tablespace OID of the relation</td>
</tr>
<tr>
<td>reldatabase</td>
<td>oid</td>
<td><code>pg_database.oid</code></td>
<td>Database OID of the relation</td>
</tr>
<tr>
<td>relforknumber</td>
<td>smallint</td>
<td></td>
<td>Fork number within the relation; see include/common/relopath.h</td>
</tr>
<tr>
<td>relblocknumber</td>
<td>bigint</td>
<td></td>
<td>Page number within the relation</td>
</tr>
<tr>
<td>isdirty</td>
<td>boolean</td>
<td></td>
<td>Is the page dirty?</td>
</tr>
</tbody>
</table>
Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>usagecount</td>
<td>smallint</td>
<td></td>
<td>Clock-sweep access count</td>
</tr>
<tr>
<td>pinning_backends</td>
<td>integer</td>
<td></td>
<td>Number of backends pinning this buffer</td>
</tr>
</tbody>
</table>

There is one row for each buffer in the shared cache. Unused buffers are shown with all fields null except bufferid. Shared system catalogs are shown as belonging to database zero.

Because the cache is shared by all the databases, there will normally be pages from relations not belonging to the current database. This means that there may not be matching join rows in `pg_class` for some rows, or that there could even be incorrect joins. If you are trying to join against `pg_class`, it’s a good idea to restrict the join to rows having `reldatabase` equal to the current database’s OID or zero.

When the `pg_buffercache` view is accessed, internal buffer manager locks are taken for long enough to copy all the buffer state data that the view will display. This ensures that the view produces a consistent set of results, while not blocking normal buffer activity longer than necessary. Nonetheless there could be some impact on database performance if this view is read often.

F.24.2. Sample Output

```sql
regression=# SELECT c.relname, count(*) AS buffers
FROM pg_buffercache b INNER JOIN pg_class c
ON b.relfilenode = pg_relation_filenode(c.oid) AND
b.reldatabase IN (0, (SELECT oid FROM pg_database
WHERE datname = current_database()))
GROUP BY c.relname
ORDER BY 2 DESC
LIMIT 10;
```

```
<table>
<thead>
<tr>
<th>relname</th>
<th>buffers</th>
</tr>
</thead>
<tbody>
<tr>
<td>tenk2</td>
<td>345</td>
</tr>
<tr>
<td>tenk1</td>
<td>141</td>
</tr>
<tr>
<td>pg_proc</td>
<td>46</td>
</tr>
<tr>
<td>pg_class</td>
<td>45</td>
</tr>
<tr>
<td>pg_attribute</td>
<td>43</td>
</tr>
<tr>
<td>pg_class_relname_nsp_index</td>
<td>30</td>
</tr>
<tr>
<td>pg_proc_proname_args_nsp_index</td>
<td>28</td>
</tr>
<tr>
<td>pg_attribute_relid_attnam_index</td>
<td>26</td>
</tr>
<tr>
<td>pg_depend</td>
<td>22</td>
</tr>
<tr>
<td>pg_depend_reference_index</td>
<td>20</td>
</tr>
</tbody>
</table>
```

(10 rows)

F.24.3. Authors

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Design suggestions: Neil Conway <neilc@samurai.com>

Debugging advice: Tom Lane <tgl@sss.pgh.pa.us>
F.25. pgcrypto

The `pgcrypto` module provides cryptographic functions for PostgreSQL.

**F.25.1. General Hashing Functions**

F.25.1.1. `digest()`

```sql
igest(data text, type text) returns bytea
igest(data bytea, type text) returns bytea
```

Computes a binary hash of the given `data`. `type` is the algorithm to use. Standard algorithms are `md5`, `sha1`, `sha224`, `sha256`, `sha384` and `sha512`. If `pgcrypto` was built with OpenSSL, more algorithms are available, as detailed in Table F-20.

If you want the digest as a hexadecimal string, use `encode()` on the result. For example:

```sql
CREATE OR REPLACE FUNCTION sha1(bytea) returns text AS $$
    SELECT encode(digest($1, 'sha1'), 'hex')
$$ LANGUAGE SQL STRICT IMMUTABLE;
```

F.25.1.2. `hmac()`

```sql
hmac(data text, key text, type text) returns bytea
hmac(data bytea, key bytea, type text) returns bytea
```

Calculates hashed MAC for `data` with key `key`. `type` is the same as in `digest()`.

This is similar to `digest()` but the hash can only be recalculated knowing the key. This prevents the scenario of someone altering data and also changing the hash to match.

If the key is larger than the hash block size it will first be hashed and the result will be used as key.

**F.25.2. Password Hashing Functions**

The functions `crypt()` and `gen_salt()` are specifically designed for hashing passwords. `crypt()` does the hashing and `gen_salt()` prepares algorithm parameters for it.

The algorithms in `crypt()` differ from the usual MD5 or SHA1 hashing algorithms in the following respects:

1. They are slow. As the amount of data is so small, this is the only way to make brute-forcing passwords hard.
2. They use a random value, called the `salt`, so that users having the same password will have different encrypted passwords. This is also an additional defense against reversing the algorithm.
3. They include the algorithm type in the result, so passwords hashed with different algorithms can co-exist.
4. Some of them are adaptive — that means when computers get faster, you can tune the algorithm to be slower, without introducing incompatibility with existing passwords.

Table F-17 lists the algorithms supported by the `crypt()` function.

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Max Password Length</th>
<th>Adaptive?</th>
<th>Salt Bits</th>
<th>Output Length</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>bf</td>
<td>72</td>
<td>yes</td>
<td>128</td>
<td>60</td>
<td>Blowfish-based, variant 2a</td>
</tr>
<tr>
<td>md5</td>
<td>unlimited</td>
<td>no</td>
<td>48</td>
<td>34</td>
<td>MD5-based crypt</td>
</tr>
<tr>
<td>xdes</td>
<td>8</td>
<td>yes</td>
<td>24</td>
<td>20</td>
<td>Extended DES</td>
</tr>
<tr>
<td>des</td>
<td>8</td>
<td>no</td>
<td>12</td>
<td>13</td>
<td>Original UNIX crypt</td>
</tr>
</tbody>
</table>

**F.25.2.1. crypt()**

`crypt(password text, salt text) returns text`

Calculates a crypt(3)-style hash of `password`. When storing a new password, you need to use `gen_salt()` to generate a new `salt` value. To check a password, pass the stored hash value as `salt`, and test whether the result matches the stored value.

Example of setting a new password:

```
UPDATE ... SET pswhash = crypt('new password', gen_salt('md5'));
```

Example of authentication:

```
SELECT (pswhash = crypt('entered password', pswhash)) AS pswmatch FROM ... ;
```

This returns `true` if the entered password is correct.

**F.25.2.2. gen_salt()**

`gen_salt(type text [, iter_count integer ]) returns text`

Generates a new random salt string for use in `crypt()`. The salt string also tells `crypt()` which algorithm to use.

The `type` parameter specifies the hashing algorithm. The accepted types are: `des`, `xdes`, `md5` and `bf`.

The `iter_count` parameter lets the user specify the iteration count, for algorithms that have one. The higher the count, the more time it takes to hash the password and therefore the more time to break it. Although with too high a count the time to calculate a hash may be several years — which is somewhat impractical. If the `iter_count` parameter is omitted, the default iteration count is used. Allowed values for `iter_count` depend on the algorithm and are shown in Table F-18.
Table F-18. Iteration Counts for crypt()

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Default</th>
<th>Min</th>
<th>Max</th>
</tr>
</thead>
<tbody>
<tr>
<td>xdes</td>
<td>725</td>
<td>1</td>
<td>16777215</td>
</tr>
<tr>
<td>bf</td>
<td>6</td>
<td>4</td>
<td>31</td>
</tr>
</tbody>
</table>

For xdes there is an additional limitation that the iteration count must be an odd number.

To pick an appropriate iteration count, consider that the original DES crypt was designed to have the speed of 4 hashes per second on the hardware of that time. Slower than 4 hashes per second would probably dampen usability. Faster than 100 hashes per second is probably too fast.

Table F-19 gives an overview of the relative slowness of different hashing algorithms. The table shows how much time it would take to try all combinations of characters in an 8-character password, assuming that the password contains either only lower case letters, or upper- and lower-case letters and numbers. In the crypt-bf entries, the number after a slash is the iter_count parameter of gen_salt.

Table F-19. Hash Algorithm Speeds

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Hashes/sec</th>
<th>For [a-z]</th>
<th>For [A-Za-z0-9]</th>
<th>Duration relative to md5 hash</th>
</tr>
</thead>
<tbody>
<tr>
<td>crypt-bf/8</td>
<td>1792</td>
<td>4 years</td>
<td>3927 years</td>
<td>100k</td>
</tr>
<tr>
<td>crypt-bf/7</td>
<td>3648</td>
<td>2 years</td>
<td>1929 years</td>
<td>50k</td>
</tr>
<tr>
<td>crypt-bf/6</td>
<td>7168</td>
<td>1 year</td>
<td>982 years</td>
<td>25k</td>
</tr>
<tr>
<td>crypt-bf/5</td>
<td>13504</td>
<td>188 days</td>
<td>521 years</td>
<td>12.5k</td>
</tr>
<tr>
<td>crypt-md5</td>
<td>171584</td>
<td>15 days</td>
<td>41 years</td>
<td>1k</td>
</tr>
<tr>
<td>crypt-des</td>
<td>23221568</td>
<td>157.5 minutes</td>
<td>108 days</td>
<td>7</td>
</tr>
<tr>
<td>sha1</td>
<td>37774272</td>
<td>90 minutes</td>
<td>68 days</td>
<td>4</td>
</tr>
<tr>
<td>md5 (hash)</td>
<td>150085504</td>
<td>22.5 minutes</td>
<td>17 days</td>
<td>1</td>
</tr>
</tbody>
</table>

Notes:

- The machine used is an Intel Mobile Core i3.
- crypt-des and crypt-md5 algorithm numbers are taken from John the Ripper v1.6.38 -test output.
- md5 hash numbers are from mdcrack 1.2.
- sha1 numbers are from lcrack-20031130-beta.
- crypt-bf numbers are taken using a simple program that loops over 1000 8-character passwords. That way I can show the speed with different numbers of iterations. For reference: john -test shows 13506 loops/sec for crypt-bf/5. (The very small difference in results is in accordance with the fact that the crypt-bf implementation in pgcrypto is the same one used in John the Ripper.)

Note that “try all combinations” is not a realistic exercise. Usually password cracking is done with the help of dictionaries, which contain both regular words and various mutations of them. So, even somewhat word-like passwords could be cracked much faster than the above numbers suggest, while a 6-character non-word-like password may escape cracking. Or not.
F.25.3. PGP Encryption Functions

The functions here implement the encryption part of the OpenPGP (RFC 4880) standard. Supported are both symmetric-key and public-key encryption.

An encrypted PGP message consists of 2 parts, or packets:

- Packet containing a session key — either symmetric-key or public-key encrypted.
- Packet containing data encrypted with the session key.

When encrypting with a symmetric key (i.e., a password):

1. The given password is hashed using a String2Key (S2K) algorithm. This is rather similar to crypt() algorithms — purposefully slow and with random salt — but it produces a full-length binary key.
2. If a separate session key is requested, a new random key will be generated. Otherwise the S2K key will be used directly as the session key.
3. If the S2K key is to be used directly, then only S2K settings will be put into the session key packet. Otherwise the session key will be encrypted with the S2K key and put into the session key packet.

When encrypting with a public key:

1. A new random session key is generated.
2. It is encrypted using the public key and put into the session key packet.

In either case the data to be encrypted is processed as follows:

1. Optional data-manipulation: compression, conversion to UTF-8, and/or conversion of line-endings.
2. The data is prefixed with a block of random bytes. This is equivalent to using a random IV.
3. An SHA1 hash of the random prefix and data is appended.
4. All this is encrypted with the session key and placed in the data packet.

F.25.3.1. pgp_sym_encrypt()

pgp_sym_encrypt(data text, psw text [, options text ]) returns bytea
pgp_sym_encrypt_bytea(data bytea, psw text [, options text ]) returns bytea

Encrypt data with a symmetric PGP key psw. The options parameter can contain option settings, as described below.

F.25.3.2. pgp_sym_decrypt()

pgp_sym_decrypt(msg bytea, psw text [, options text ]) returns text
pgp_sym_decrypt_bytea(msg bytea, psw text [, options text ]) returns bytea

Decrypt a symmetric-key-encrypted PGP message.

Decrypting bytea data with pgp_sym_decrypt is disallowed. This is to avoid outputting invalid character data. Decrypting originally textual data with pgp_sym_decrypt_bytea is fine.
The `options` parameter can contain option settings, as described below.

**F.25.3.3. pgp_pub_encrypt()**

```sql
gpg_pub_encrypt(data text, key bytea [, options text ]) returns bytea
gpg_pub_encrypt_bytea(data bytea, key bytea [, options text ]) returns bytea
```

Encrypt `data` with a public PGP key. Giving this function a secret key will produce an error.

The `options` parameter can contain option settings, as described below.

**F.25.3.4. pgp_pub_decrypt()**

```sql
gpg_pub_decrypt(msg bytea, key bytea [, psw text [, options text ]]) returns text
gpg_pub_decrypt_bytea(msg bytea, key bytea [, psw text [, options text ]]) returns bytea
```

Decrypt a public-key-encrypted message. `key` must be the secret key corresponding to the public key that was used to encrypt. If the secret key is password-protected, you must give the password in `psw`. If there is no password, but you want to specify options, you need to give an empty password.

Decrypting `bytea` data with `pgp_pub_decrypt` is disallowed. This is to avoid outputting invalid character data. Decrypting originally textual data with `pgp_pub_decrypt_bytea` is fine.

The `options` parameter can contain option settings, as described below.

**F.25.3.5. pgp_key_id()**

```sql
gpg_key_id(bytea) returns text
```

`gpg_key_id` extracts the key ID of a PGP public or secret key. Or it gives the key ID that was used for encrypting the data, if given an encrypted message.

It can return 2 special key IDs:

- **SYMKEY**
  The message is encrypted with a symmetric key.
- **ANYKEY**
  The message is public-key encrypted, but the key ID has been removed. That means you will need to try all your secret keys on it to see which one decrypts it. *pgcrypto* itself does not produce such messages.

Note that different keys may have the same ID. This is rare but a normal event. The client application should then try to decrypt with each one, to see which fits — like handling `ANYKEY`.

**F.25.3.6. armor(), dearmor()**

```sql
armor(data bytea [, , keys text[], values text[] ]) returns text
dearmor(data text) returns bytea
```

These functions wrap/unwrap binary data into PGP ASCII-armor format, which is basically Base64 with CRC and additional formatting.
If the keys and values arrays are specified, an armor header is added to the armored format for each key/value pair. Both arrays must be single-dimensional, and they must be of the same length. The keys and values cannot contain any non-ASCII characters.

**F.25.3.7. pgp_armor_headers**

pgp_armor_headers(data text, key out text, value out text) returns setof record

pgp_armor_headers() extracts the armor headers from data. The return value is a set of rows with two columns, key and value. If the keys or values contain any non-ASCII characters, they are treated as UTF-8.

**F.25.3.8. Options for PGP Functions**

Options are named to be similar to GnuPG. An option’s value should be given after an equal sign; separate options from each other with commas. For example:

```
pgp_sym_encrypt(data, psw, 'compress-algo=1, cipher-algo=aes256')
```

All of the options except convert-crlf apply only to encrypt functions. Decrypt functions get the parameters from the PGP data.

The most interesting options are probably compress-algo and unicode-mode. The rest should have reasonable defaults.

**F.25.3.8.1. cipher-algo**

Which cipher algorithm to use.

Values: bf, aes128, aes192, aes256 (OpenSSL-only: 3des, cast5)
Default: aes128
Applies to: pgp_sym_encrypt, pgp_pub_encrypt

**F.25.3.8.2. compress-algo**

Which compression algorithm to use. Only available if PostgreSQL was built with zlib.

Values:
  0 - no compression
  1 - ZIP compression
  2 - ZLIB compression (= ZIP plus meta-data and block CRCs)
Default: 0
Applies to: pgp_sym_encrypt, pgp_pub_encrypt

**F.25.3.8.3. compress-level**

How much to compress. Higher levels compress smaller but are slower. 0 disables compression.

Values: 0, 1-9
Default: 6
Appendix F. Additional Supplied Modules

Applies to: pgp_sym_encrypt, pgp_pub_encrypt

F.25.3.8.4. convert-crlf
Whether to convert \n into \r\n when encrypting and \r\n to \n when decrypting. RFC 4880 specifies that text data should be stored using \r\n line-feeds. Use this to get fully RFC-compliant behavior.

Values: 0, 1
Default: 0
Applies to: pgp_sym_encrypt, pgp_pub_encrypt, pgp_sym_decrypt, pgp_pub_decrypt

F.25.3.8.5. disable-mdc
Do not protect data with SHA-1. The only good reason to use this option is to achieve compatibility with ancient PGP products, predating the addition of SHA-1 protected packets to RFC 4880. Recent gnupg.org and pgp.com software supports it fine.

Values: 0, 1
Default: 0
Applies to: pgp_sym_encrypt, pgp_pub_encrypt

F.25.3.8.6. sess-key
Use separate session key. Public-key encryption always uses a separate session key; this option is for symmetric-key encryption, which by default uses the S2K key directly.

Values: 0, 1
Default: 0
Applies to: pgp_sym_encrypt

F.25.3.8.7. s2k-mode
Which S2K algorithm to use.

Values:
   0 - Without salt. Dangerous!
   1 - With salt but with fixed iteration count.
   3 - Variable iteration count.
Default: 3
Applies to: pgp_sym_encrypt

F.25.3.8.8. s2k-count
The number of iterations of the S2K algorithm to use. It must be a value between 1024 and 65011712, inclusive.

Default: A random value between 65536 and 253952
Applies to: pgp_sym_encrypt, only with s2k-mode=3
Appendix F. Additional Supplied Modules

**F.25.3.8.9. s2k-digest-algo**

Which digest algorithm to use in S2K calculation.

Values: md5, sha1
Default: sha1
Applies to: pgp_sym_encrypt

**F.25.3.8.10. s2k-cipher-algo**

Which cipher to use for encrypting separate session key.

Values: bf, aes, aes128, aes192, aes256
Default: use cipher-algo
Applies to: pgp_sym_encrypt

**F.25.3.8.11. unicode-mode**

Whether to convert textual data from database internal encoding to UTF-8 and back. If your database already is UTF-8, no conversion will be done, but the message will be tagged as UTF-8. Without this option it will not be.

Values: 0, 1
Default: 0
Applies to: pgp_sym_encrypt, pgp_pub_encrypt

**F.25.3.9. Generating PGP Keys with GnuPG**

To generate a new key:

```
gpg --gen-key
```

The preferred key type is “DSA and Elgamal”.

For RSA encryption you must create either DSA or RSA sign-only key as master and then add an RSA encryption subkey with `gpg --edit-key`.

To list keys:

```
gpg --list-secret-keys
```

To export a public key in ASCII-armor format:

```
gpg -a --export KEYID > public.key
```

To export a secret key in ASCII-armor format:

```
gpg -a --export-secret-keys KEYID > secret.key
```
Appendix F. Additional Supplied Modules

You need to use `dearmor()` on these keys before giving them to the PGP functions. Or if you can handle binary data, you can drop `-a` from the command.

For more details see `man gpg`, The GNU Privacy Handbook\(^4\) and other documentation on http://www.gnupg.org.

F.25.3.10. Limitations of PGP Code

- No support for signing. That also means that it is not checked whether the encryption subkey belongs to the master key.
- No support for encryption key as master key. As such practice is generally discouraged, this should not be a problem.
- No support for several subkeys. This may seem like a problem, as this is common practice. On the other hand, you should not use your regular GPG/PGP keys with `pgcrypto`, but create new ones, as the usage scenario is rather different.

F.25.4. Raw Encryption Functions

These functions only run a cipher over data; they don’t have any advanced features of PGP encryption. Therefore they have some major problems:

1. They use user key directly as cipher key.
2. They don’t provide any integrity checking, to see if the encrypted data was modified.
3. They expect that users manage all encryption parameters themselves, even IV.
4. They don’t handle text.

So, with the introduction of PGP encryption, usage of raw encryption functions is discouraged.

```sql
encrypt(data bytea, key bytea, type text) returns bytea
decrypt(data bytea, key bytea, type text) returns bytea
encrypt_iv(data bytea, key bytea, iv bytea, type text) returns bytea
decrypt_iv(data bytea, key bytea, iv bytea, type text) returns bytea
```

Encrypt/decrypt data using the cipher method specified by `type`. The syntax of the `type` string is:

```sql
algorithm [ - mode ] [ /pad: padding ]
```

where `algorithm` is one of:

- `bf` — Blowfish
- `aes` — AES (Rijndael-128, -192 or -256)

and `mode` is one of:

---

Appendix F. Additional Supplied Modules

• cbc — next block depends on previous (default)
• ecb — each block is encrypted separately (for testing only)

and padding is one of:

• pkcs — data may be any length (default)
• none — data must be multiple of cipher block size

So, for example, these are equivalent:

```sql
encrypt(data, 'fooz', 'bf')
encrypt(data, 'fooz', 'bf-cbc/pad:pkcs')
```

In `encrypt_iv` and `decrypt_iv`, the iv parameter is the initial value for the CBC mode; it is ignored for ECB. It is clipped or padded with zeroes if not exactly block size. It defaults to all zeroes in the functions without this parameter.

F.25.5. Random-Data Functions

`gen_random_bytes(count integer) returns bytea`

Returns `count` cryptographically strong random bytes. At most 1024 bytes can be extracted at a time. This is to avoid draining the randomness generator pool.

`gen_random_uuid() returns uuid`

Returns a version 4 (random) UUID.

F.25.6. Notes

F.25.6.1. Configuration

`pgcrypto` configures itself according to the findings of the main `PostgreSQL configure` script. The options that affect it are `--with-zlib` and `--with-openssl`.

When compiled with zlib, PGP encryption functions are able to compress data before encrypting.

When compiled with OpenSSL, there will be more algorithms available. Also public-key encryption functions will be faster as OpenSSL has more optimized `BIGNUM` functions.

Table F-20. Summary of Functionality with and without OpenSSL

<table>
<thead>
<tr>
<th>Functionality</th>
<th>Built-in</th>
<th>With OpenSSL</th>
</tr>
</thead>
<tbody>
<tr>
<td>MD5</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>SHA1</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>SHA224/256/384/512</td>
<td>yes</td>
<td>yes (Note 1)</td>
</tr>
<tr>
<td>Other digest algorithms</td>
<td>no</td>
<td>yes (Note 2)</td>
</tr>
<tr>
<td>Blowfish</td>
<td>yes</td>
<td>yes</td>
</tr>
</tbody>
</table>
Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Functionality</th>
<th>Built-in</th>
<th>With OpenSSL</th>
</tr>
</thead>
<tbody>
<tr>
<td>AES</td>
<td>yes</td>
<td>yes (Note 3)</td>
</tr>
<tr>
<td>DES/3DES/CAST5</td>
<td>no</td>
<td>yes</td>
</tr>
<tr>
<td>Raw encryption</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>PGP Symmetric encryption</td>
<td>yes</td>
<td>yes</td>
</tr>
<tr>
<td>PGP Public-Key encryption</td>
<td>yes</td>
<td>yes</td>
</tr>
</tbody>
</table>

Notes:

1. SHA2 algorithms were added to OpenSSL in version 0.9.8. For older versions, `pgcrypto` will use built-in code.
2. Any digest algorithm OpenSSL supports is automatically picked up. This is not possible with ciphers, which need to be supported explicitly.
3. AES is included in OpenSSL since version 0.9.7. For older versions, `pgcrypto` will use built-in code.

F.25.6.2. NULL Handling

As is standard in SQL, all functions return NULL, if any of the arguments are NULL. This may create security risks on careless usage.

F.25.6.3. Security Limitations

All `pgcrypto` functions run inside the database server. That means that all the data and passwords move between `pgcrypto` and client applications in clear text. Thus you must:

1. Connect locally or use SSL connections.
2. Trust both system and database administrator.

If you cannot, then better do crypto inside client application.

The implementation does not resist side-channel attacks\(^5\). For example, the time required for a `pgcrypto` decryption function to complete varies among ciphertexts of a given size.

F.25.6.4. Useful Reading

  The GNU Privacy Handbook.
  Describes the crypt-blowfish algorithm.
  How to choose a good password.

Interesting idea for picking passwords.
- [http://www.interhack.net/people/cmcurtin/snake-oil-faq.html](http://www.interhack.net/people/cmcurtin/snake-oil-faq.html)
  Describes good and bad cryptography.

**F.25.6.5. Technical References**

  OpenPGP message format.
  The MD5 Message-Digest Algorithm.
  HMAC: Keyed-Hashing for Message Authentication.
  Comparison of crypt-des, crypt-md5 and bcrypt algorithms.
  Description of Fortuna CSPRNG.
- [http://jlcooke.ca/random/](http://jlcooke.ca/random/)
  Jean-Luc Cooke Fortuna-based /dev/random driver for Linux.

**F.25.7. Author**

Marko Kreen <markokr@gmail.com>

`pgcrypto` uses code from the following sources:

<table>
<thead>
<tr>
<th>Algorithm</th>
<th>Author</th>
<th>Source origin</th>
</tr>
</thead>
<tbody>
<tr>
<td>DES crypt</td>
<td>David Burren and others</td>
<td>FreeBSD libcrypt</td>
</tr>
<tr>
<td>MD5 crypt</td>
<td>Poul-Henning Kamp</td>
<td>FreeBSD libcrypt</td>
</tr>
<tr>
<td>Blowfish crypt</td>
<td>Solar Designer</td>
<td><a href="http://www.openwall.com">www.openwall.com</a></td>
</tr>
<tr>
<td>Blowfish cipher</td>
<td>Simon Tatham</td>
<td>PuTTY</td>
</tr>
<tr>
<td>Rijndael cipher</td>
<td>Brian Gladman</td>
<td>OpenBSD sys/crypto</td>
</tr>
<tr>
<td>MD5 hash and SHA1</td>
<td>WIDE Project</td>
<td>KAME kame/sys/crypto</td>
</tr>
<tr>
<td>SHA256/384/512</td>
<td>Aaron D. Gifford</td>
<td>OpenBSD sys/crypto</td>
</tr>
<tr>
<td>BIGNUM math</td>
<td>Michael J. Fromberger</td>
<td>dartmouth.edu/~sting/sw/imath</td>
</tr>
</tbody>
</table>

**F.26. pg_freespacemap**

The `pg_freespacemap` module provides a means for examining the free space map (FSM). It provides a function called `pg_freespace`, or two overloaded functions, to be precise. The functions
show the value recorded in the free space map for a given page, or for all pages in the relation.
By default public access is revoked from the functions, just in case there are security issues lurking.

F.26.1. Functions

pg_freespace(rel regclass IN, blkno bigint IN) returns int2

Returns the amount of free space on the page of the relation, specified by blkno, according to the FSM.

pg_freespace(rel regclass IN, blkno OUT bigint, avail OUT int2)

Displays the amount of free space on each page of the relation, according to the FSM. A set of (blkno bigint, avail int2) tuples is returned, one tuple for each page in the relation.

The values stored in the free space map are not exact. They’re rounded to precision of 1/256th of BLCKSZ (32 bytes with default BLCKSZ), and they’re not kept fully up-to-date as tuples are inserted and updated.

For indexes, what is tracked is entirely-unused pages, rather than free space within pages. Therefore, the values are not meaningful, just whether a page is full or empty.

Note: The interface was changed in version 8.4, to reflect the new FSM implementation introduced in the same version.

F.26.2. Sample Output

postgres=# SELECT * FROM pg_freespace('foo');

| blkno | avail |
|-------+-------|
| 0     |  0    |
| 1     |  0    |
| 2     |  32   |
| 3     |  704  |
| 4     |  704  |
| 5     |  704  |
| 6     |  704  |
| 7     | 1216  |
| 8     |  704  |
| 9     |  704  |
| 10    |  704  |
| 11    |  704  |
| 12    |  704  |
| 13    |  704  |
| 14    |  704  |
| 15    |  704  |
| 16    |  704  |
| 17    |  704  |
| 18    |  704  |
| 19    | 3648  |

(20 rows)

postgres=# SELECT * FROM pg_freespace('foo', 7);
Appendix F. Additional Supplied Modules

F.26.3. Author

Original version by Mark Kirkwood <markir@paradise.net.nz>. Rewritten in version 8.4 to suit new FSM implementation by Heikki Linnakangas <heikki@enterprisedb.com>

F.27. pg_prewarm

The pg_prewarm module provides a convenient way to load relation data into either the operating system buffer cache or the PostgreSQL buffer cache.

F.27.1. Functions

pg_prewarm(regclass, mode text default 'buffer', fork text default 'main',
            first_block int8 default null,
            last_block int8 default null) RETURNS int8

The first argument is the relation to be prewarmed. The second argument is the prewarming method to be used, as further discussed below; the third is the relation fork to be prewarmed, usually main. The fourth argument is the first block number to prewarm (NULL is accepted as a synonym for zero). The fifth argument is the last block number to prewarm (NULL means prewarm through the last block in the relation). The return value is the number of blocks prewarmed.

There are three available prewarming methods. prefetch issues asynchronous prefetch requests to the operating system, if this is supported, or throws an error otherwise. read reads the requested range of blocks; unlike prefetch, this is synchronous and supported on all platforms and builds, but may be slower. buffer reads the requested range of blocks into the database buffer cache.

Note that with any of these methods, attempting to prewarm more blocks than can be cached — by the OS when using prefetch or read, or by PostgreSQL when using buffer — will likely result in lower-numbered blocks being evicted as higher numbered blocks are read in. Prewarmed data also enjoys no special protection from cache evictions, so it is possible that other system activity may evict the newly prewarmed blocks shortly after they are read; conversely, prewarming may also evict other data from cache. For these reasons, prewarming is typically most useful at startup, when caches are largely empty.

F.27.2. Author

Robert Haas <rhaas@postgresql.org>
Appendix F. Additional Supplied Modules

F.28. pgrowlocks

The pgrowlocks module provides a function to show row locking information for a specified table.

F.28.1. Overview

pgrowlocks(text) returns setof record

The parameter is the name of a table. The result is a set of records, with one row for each locked row within the table. The output columns are shown in Table F-21.

Table F-21. pgrowlocks Output Columns

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>locked_row</td>
<td>tid</td>
<td>Tuple ID (TID) of locked row</td>
</tr>
<tr>
<td>locker</td>
<td>xid</td>
<td>Transaction ID of locker, or multiact ID if multitransaction</td>
</tr>
<tr>
<td>multi</td>
<td>boolean</td>
<td>True if locker is a multitransaction</td>
</tr>
<tr>
<td>xids</td>
<td>xid[]</td>
<td>Transaction IDs of lockers (more than one if multitransaction)</td>
</tr>
<tr>
<td>modes</td>
<td>text[]</td>
<td>Lock mode of lockers (more than one if multitransaction), an array of Key Share, Share, For No Key Update, No Key Update, For Update, Update.</td>
</tr>
<tr>
<td>pids</td>
<td>integer[]</td>
<td>Process IDs of locking backends (more than one if multitransaction)</td>
</tr>
</tbody>
</table>

pgrowlocks takes AccessShareLock for the target table and reads each row one by one to collect the row locking information. This is not very speedy for a large table. Note that:

1. If the table as a whole is exclusive-locked by someone else, pgrowlocks will be blocked.
2. pgrowlocks is not guaranteed to produce a self-consistent snapshot. It is possible that a new row lock is taken, or an old lock is freed, during its execution.

pgrowlocks does not show the contents of locked rows. If you want to take a look at the row contents at the same time, you could do something like this:

```
SELECT * FROM accounts AS a, pgrowlocks('accounts') AS p
WHERE p.locked_row = a.ctid;
```

Be aware however that such a query will be very inefficient.
F.28.2. Sample Output

```sql
-\# SELECT * FROM pgrowlocks('t1');

<table>
<thead>
<tr>
<th>locked_row</th>
<th>locker</th>
<th>multi</th>
<th>xids</th>
<th>modes</th>
<th>pids</th>
</tr>
</thead>
<tbody>
<tr>
<td>(0,1)</td>
<td>609</td>
<td>f</td>
<td>{609}</td>
<td>&quot;For Share&quot;</td>
<td>{3161}</td>
</tr>
<tr>
<td>(0,2)</td>
<td>609</td>
<td>f</td>
<td>{609}</td>
<td>&quot;For Share&quot;</td>
<td>{3161}</td>
</tr>
<tr>
<td>(0,3)</td>
<td>607</td>
<td>f</td>
<td>{607}</td>
<td>&quot;For Update&quot;</td>
<td>{3107}</td>
</tr>
<tr>
<td>(0,4)</td>
<td>607</td>
<td>f</td>
<td>{607}</td>
<td>&quot;For Update&quot;</td>
<td>{3107}</td>
</tr>
</tbody>
</table>
```

(4 rows)

F.28.3. Author

Tatsuo Ishii

F.29. pg_stat_statements

The `pg_stat_statements` module provides a means for tracking execution statistics of all SQL statements executed by a server.

The module must be loaded by adding `pg_stat_statements` to `shared_preload_libraries` in `postgresql.conf`, because it requires additional shared memory. This means that a server restart is needed to add or remove the module.

When `pg_stat_statements` is loaded, it tracks statistics across all databases of the server. To access and manipulate these statistics, the module provides a view, `pg_stat_statements`, and the utility functions `pg_stat_statements_reset` and `pg_stat_statements`. These are not available globally but can be enabled for a specific database with `CREATE EXTENSION pg_stat_statements`.

F.29.1. The `pg_stat_statements` View

The statistics gathered by the module are made available via a view named `pg_stat_statements`. This view contains one row for each distinct database ID, user ID and query ID (up to the maximum number of distinct statements that the module can track). The columns of the view are shown in Table F-22.

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>userid</td>
<td>oid</td>
<td>pg_authid.oid</td>
<td>OID of user who executed the statement</td>
</tr>
<tr>
<td>dbid</td>
<td>oid</td>
<td>pg_database.oid</td>
<td>OID of database in which the statement was executed</td>
</tr>
<tr>
<td>queryid</td>
<td>bigint</td>
<td></td>
<td>Internal hash code, computed from the statement’s parse tree</td>
</tr>
<tr>
<td>Name</td>
<td>Type</td>
<td>References</td>
<td>Description</td>
</tr>
<tr>
<td>---------------</td>
<td>-----------------</td>
<td>------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>query</td>
<td>text</td>
<td></td>
<td>Text of a representative statement</td>
</tr>
<tr>
<td>calls</td>
<td>bigint</td>
<td></td>
<td>Number of times executed</td>
</tr>
<tr>
<td>total_time</td>
<td>double precision</td>
<td></td>
<td>Total time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>min_time</td>
<td>double precision</td>
<td></td>
<td>Minimum time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>max_time</td>
<td>double precision</td>
<td></td>
<td>Maximum time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>mean_time</td>
<td>double precision</td>
<td></td>
<td>Mean time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>stddev_time</td>
<td>double precision</td>
<td></td>
<td>Population standard deviation of time spent in the statement, in milliseconds</td>
</tr>
<tr>
<td>rows</td>
<td>bigint</td>
<td></td>
<td>Total number of rows retrieved or affected by the statement</td>
</tr>
<tr>
<td>shared_blks_hit</td>
<td>bigint</td>
<td></td>
<td>Total number of shared block cache hits by the statement</td>
</tr>
<tr>
<td>shared_blks_read</td>
<td>bigint</td>
<td></td>
<td>Total number of shared blocks read by the statement</td>
</tr>
<tr>
<td>shared_blks_dirtied</td>
<td>bigint</td>
<td></td>
<td>Total number of shared blocks dirtied by the statement</td>
</tr>
<tr>
<td>shared_blks_written</td>
<td>bigint</td>
<td></td>
<td>Total number of shared blocks written by the statement</td>
</tr>
<tr>
<td>local_blks_hit</td>
<td>bigint</td>
<td></td>
<td>Total number of local block cache hits by the statement</td>
</tr>
<tr>
<td>local_blks_read</td>
<td>bigint</td>
<td></td>
<td>Total number of local blocks read by the statement</td>
</tr>
<tr>
<td>local_blks_dirtied</td>
<td>bigint</td>
<td></td>
<td>Total number of local blocks dirtied by the statement</td>
</tr>
<tr>
<td>local_blks_written</td>
<td>bigint</td>
<td></td>
<td>Total number of local blocks written by the statement</td>
</tr>
</tbody>
</table>
Appendix F: Additional Supplied Modules

<table>
<thead>
<tr>
<th>Name</th>
<th>Type</th>
<th>References</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>temp_blks_read</td>
<td>bigint</td>
<td></td>
<td>Total number of temp blocks read by the statement</td>
</tr>
<tr>
<td>temp_blks_written</td>
<td>bigint</td>
<td></td>
<td>Total number of temp blocks written by the statement</td>
</tr>
<tr>
<td>blk_read_time</td>
<td>double precision</td>
<td></td>
<td>Total time the statement spent reading blocks, in milliseconds (if track_io_timing is enabled, otherwise zero)</td>
</tr>
<tr>
<td>blk_write_time</td>
<td>double precision</td>
<td></td>
<td>Total time the statement spent writing blocks, in milliseconds (if track_io_timing is enabled, otherwise zero)</td>
</tr>
</tbody>
</table>

For security reasons, non-superusers are not allowed to see the SQL text or queryid of queries executed by other users. They can see the statistics, however, if the view has been installed in their database.

Plannable queries (that is, SELECT, INSERT, UPDATE, and DELETE) are combined into a single pg_stat_statements entry whenever they have identical query structures according to an internal hash calculation. Typically, two queries will be considered the same for this purpose if they are semantically equivalent except for the values of literal constants appearing in the query. Utility commands (that is, all other commands) are compared strictly on the basis of their textual query strings, however.

When a constant’s value has been ignored for purposes of matching the query to other queries, the constant is replaced by ? in the pg_stat_statements display. The rest of the query text is that of the first query that had the particular queryid hash value associated with the pg_stat_statements entry.

In some cases, queries with visibly different texts might get merged into a single pg_stat_statements entry. Normally this will happen only for semantically equivalent queries, but there is a small chance of hash collisions causing unrelated queries to be merged into one entry. (This cannot happen for queries belonging to different users or databases, however.)

Since the queryid hash value is computed on the post-parse-analysis representation of the queries, the opposite is also possible: queries with identical texts might appear as separate entries, if they have different meanings as a result of factors such as different search_path settings.

Consumers of pg_stat_statements may wish to use queryid (perhaps in combination with dbid and userid) as a more stable and reliable identifier for each entry than its query text. However, it is important to understand that there are only limited guarantees around the stability of the queryid hash value. Since the identifier is derived from the post-parse-analysis tree, its value is a function of, among other things, the internal object identifiers appearing in this representation. This has some counterintuitive implications. For example, pg_stat_statements will consider two apparently-identical queries to be distinct, if they reference a table that was dropped and recreated between the executions of the two queries. The hashing process is also sensitive to differences in machine architecture and other facets of the platform. Furthermore, it is not safe to assume that queryid will...
be stable across major versions of PostgreSQL.

As a rule of thumb, \texttt{queryid} values can be assumed to be stable and comparable only so long as the underlying server version and catalog metadata details stay exactly the same. Two servers participating in replication based on physical WAL replay can be expected to have identical \texttt{queryid} values for the same query. However, logical replication schemes do not promise to keep replicas identical in all relevant details, so \texttt{queryid} will not be a useful identifier for accumulating costs across a set of logical replicas. If in doubt, direct testing is recommended.

The representative query texts are kept in an external disk file, and do not consume shared memory. Therefore, even very lengthy query texts can be stored successfully. However, if many long query texts are accumulated, the external file might grow unmanageably large. As a recovery method if that happens, \texttt{pg_stat_statements} may choose to discard the query texts, whereupon all existing entries in the \texttt{pg_stat_statements} view will show null query fields, though the statistics associated with each \texttt{queryid} are preserved. If this happens, consider reducing \texttt{pg_stat_statements.max} to prevent recurrences.

\section*{F.29.2. Functions}

\texttt{pg_stat_statements_reset()} returns void

\texttt{pg_stat_statements_reset} discards all statistics gathered so far by \texttt{pg_stat_statements}. By default, this function can only be executed by superusers.

\texttt{pg_stat_statements(showtext boolean)} returns setof record

The \texttt{pg_stat_statements} view is defined in terms of a function also named \texttt{pg_stat_statements}. It is possible for clients to call the \texttt{pg_stat_statements} function directly, and by specifying \texttt{showtext := false} have query text be omitted (that is, the OUT argument that corresponds to the view’s query column will return nulls). This feature is intended to support external tools that might wish to avoid the overhead of repeatedly retrieving query texts of indeterminate length. Such tools can instead cache the first query text observed for each entry themselves, since that is all \texttt{pg_stat_statements} itself does, and then retrieve query texts only as needed. Since the server stores query texts in a file, this approach may reduce physical I/O for repeated examination of the \texttt{pg_stat_statements} data.

\section*{F.29.3. Configuration Parameters}

\texttt{pg_stat_statements.max(integer)}

\texttt{pg_stat_statements.max} is the maximum number of statements tracked by the module (i.e., the maximum number of rows in the \texttt{pg_stat_statements} view). If more distinct statements than that are observed, information about the least-executed statements is discarded. The default value is 5000. This parameter can only be set at server start.

\texttt{pg_stat_statements.track(enum)}

\texttt{pg_stat_statements.track} controls which statements are counted by the module. Specify \texttt{top} to track top-level statements (those issued directly by clients), \texttt{all} to also track nested statements (such as statements invoked within functions), or \texttt{none} to disable statement statistics collection. The default value is \texttt{top}. Only superusers can change this setting.
Appendix F. Additional Supplied Modules

pg_stat_statements.track_utility (boolean)

pg_stat_statements.track_utility controls whether utility commands are tracked by the module. Utility commands are all those other than SELECT, INSERT, UPDATE and DELETE. The default value is on. Only superusers can change this setting.

pg_stat_statements.save (boolean)

pg_stat_statements.save specifies whether to save statement statistics across server shut-downs. If it is off then statistics are not saved at shutdown nor reloaded at server start. The default value is on. This parameter can only be set in the postgresql.conf file or on the server command line.

The module requires additional shared memory proportional to pg_stat_statements.max. Note that this memory is consumed whenever the module is loaded, even if pg_stat_statements.track is set to none.

These parameters must be set in postgresql.conf. Typical usage might be:

```
# postgresql.conf
shared_preload_libraries = 'pg_stat_statements'

pg_stat_statements.max = 10000
pg_stat_statements.track = all
```

F.29.4. Sample Output

```
bench=# SELECT pg_stat_statements_reset();

$ pgbench -i bench
$ pgbench -c10 -t300 bench

bench=# \x

bench=# SELECT query, calls, total_time, rows, 100.0 * shared_blks_hit / nullif(shared_blks_hit + shared_blks_read, 0) AS hit_percent
FROM pg_stat_statements ORDER BY total_time DESC LIMIT 5;

- [ RECORD 1 ]---------------------------------------------------------------
| query | UPDATE pgbench_branches SET bbalance = bbalance + ? WHERE bid = ?; |
| calls | 3000 |
| total_time | 9609.00100000002 |
| rows | 2836 |
| hit_percent | 99.9778970000200936 |

- [ RECORD 2 ]---------------------------------------------------------------
| query | UPDATE pgbench_tellers SET tbalance = tbalance + ? WHERE tid = ?; |
| calls | 3000 |
| total_time | 8015.156 |
| rows | 2990 |
| hit_percent | 99.9731126579631345 |

- [ RECORD 3 ]---------------------------------------------------------------
| query | copy pgbench_accounts from stdin |
| calls | 1 |
| total_time | 310.624 |
| rows | 100000 |
| hit_percent | 0.30395136778115501520 |

- [ RECORD 4 ]---------------------------------------------------------------
```
Appendix F. Additional Supplied Modules

```
query | UPDATE pgbench_accounts SET abalance = abalance + ? WHERE aid = ?;
calls | 3000
total_time | 271.74199999997997
rows | 3000
hit_percent | 93.7968855088209426
```

F.29.5. Authors

Takahiro Itagaki <itagaki.takahiro@oss.ntt.co.jp>. Query normalization added by Peter Geoghegan <peter@2ndquadrant.com>.

F.30. pgstattuple

The pgstattuple module provides various functions to obtain tuple-level statistics.

F.30.1. Functions

```
pgstattuple(regclass) returns record
```

pgstattuple returns a relation’s physical length, percentage of “dead” tuples, and other info. This may help users to determine whether vacuum is necessary or not. The argument is the target relation’s name (optionally schema-qualified) or OID. For example:

```
test=> SELECT * FROM pgstattuple('pg_catalog.pg_proc');
```

```
- [ RECORD 1 ]-------+
  table_len | 458752
  tuple_count | 1470
  tuple_len | 438896
  tuple_percent | 95.67
  dead_tuple_count | 11
  dead_tuple_len | 3157
  dead_tuple_percent | 0.69
  free_space | 8932
  free_percent | 1.95
```

The output columns are described in Table F-23.

Table F-23. pgstattuple Output Columns

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_len</td>
<td>bigint</td>
<td>Physical relation length in bytes</td>
</tr>
<tr>
<td>tuple_count</td>
<td>bigint</td>
<td>Number of live tuples</td>
</tr>
<tr>
<td>tuple_len</td>
<td>bigint</td>
<td>Total length of live tuples in bytes</td>
</tr>
<tr>
<td>Column</td>
<td>Type</td>
<td>Description</td>
</tr>
<tr>
<td>--------------------</td>
<td>----------</td>
<td>------------------------------------</td>
</tr>
<tr>
<td>tuple_percent</td>
<td>float8</td>
<td>Percentage of live tuples</td>
</tr>
<tr>
<td>dead_tuple_count</td>
<td>bigint</td>
<td>Number of dead tuples</td>
</tr>
<tr>
<td>dead_tuple_len</td>
<td>bigint</td>
<td>Total length of dead tuples in bytes</td>
</tr>
<tr>
<td>dead_tuple_percent</td>
<td>float8</td>
<td>Percentage of dead tuples</td>
</tr>
<tr>
<td>free_space</td>
<td>bigint</td>
<td>Total free space in bytes</td>
</tr>
<tr>
<td>free_percent</td>
<td>float8</td>
<td>Percentage of free space</td>
</tr>
</tbody>
</table>

Note: The table_len will always be greater than the sum of the tuple_len, dead_tuple_len and free_space. The difference is accounted for by fixed page overhead, the per-page table of pointers to tuples, and padding to ensure that tuples are correctly aligned.

`pgstattuple` acquires only a read lock on the relation. So the results do not reflect an instantaneous snapshot; concurrent updates will affect them.

`pgstattuple` judges a tuple is “dead” if `HeapTupleSatisfiesDirty` returns false.

`pgstattuple(text)` returns record

This is the same as `pgstattuple(regclass)`, except that the target relation is specified as TEXT. This function is kept because of backward-compatibility so far, and will be deprecated in some future release.

`pgstatindex(regclass)` returns record

`pgstatindex` returns a record showing information about a B-tree index. For example:

```sql
test=> SELECT * FROM pgstatindex('pg_cast_oid_index');
- RECORD 1 ---------
| version | 2 |
| tree_level | 0 |
| index_size | 16384 |
| root_block_no | 1 |
| internal_pages | 0 |
| leaf_pages | 1 |
| empty_pages | 0 |
| deleted_pages | 0 |
| avg_leaf_density | 54.27 |
| leaf_fragmentation | 0 |
```

The output columns are:

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>version</td>
<td>integer</td>
<td>B-tree version number</td>
</tr>
<tr>
<td>tree_level</td>
<td>integer</td>
<td>Tree level of the root page</td>
</tr>
<tr>
<td>index_size</td>
<td>bigint</td>
<td>Total index size in bytes</td>
</tr>
<tr>
<td>root_block_no</td>
<td>bigint</td>
<td>Location of root page (zero if none)</td>
</tr>
<tr>
<td>internal_pages</td>
<td>bigint</td>
<td>Number of “internal” (upper-level) pages</td>
</tr>
<tr>
<td>leaf_pages</td>
<td>bigint</td>
<td>Number of leaf pages</td>
</tr>
</tbody>
</table>
Appendix F: Additional Supplied Modules

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>empty_pages</td>
<td>bigint</td>
<td>Number of empty pages</td>
</tr>
<tr>
<td>deleted_pages</td>
<td>bigint</td>
<td>Number of deleted pages</td>
</tr>
<tr>
<td>avg_leaf_density</td>
<td>float8</td>
<td>Average density of leaf pages</td>
</tr>
<tr>
<td>leaf_fragmentation</td>
<td>float8</td>
<td>Leaf page fragmentation</td>
</tr>
</tbody>
</table>

The reported `index_size` will normally correspond to one more page than is accounted for by `internal_pages + leaf_pages + empty_pages + deleted_pages`, because it also includes the index’s metapage.

As with `pgstattuple`, the results are accumulated page-by-page, and should not be expected to represent an instantaneous snapshot of the whole index.

`pgstatindex(text)` returns record

This is the same as `pgstatindex(regclass)`, except that the target index is specified as TEXT. This function is kept because of backward-compatibility so far, and will be deprecated in some future release.

`pgstatginindex(regclass)` returns record

`pgstatginindex` returns a record showing information about a GIN index. For example:

```sql
test=> SELECT * FROM pgstatginindex('test_gin_index');
- [ RECORD 1 ]-----
version | 1
pending_pages | 0
pending_tuples | 0
```

The output columns are:

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>version</td>
<td>integer</td>
<td>GIN version number</td>
</tr>
<tr>
<td>pending_pages</td>
<td>integer</td>
<td>Number of pages in the pending list</td>
</tr>
<tr>
<td>pending_tuples</td>
<td>bigint</td>
<td>Number of tuples in the pending list</td>
</tr>
</tbody>
</table>

`pg_relpages(regclass)` returns bigint

`pg_relpages` returns the number of pages in the relation.

`pg_relpages(text)` returns bigint

This is the same as `pg_relpages(regclass)`, except that the target relation is specified as TEXT. This function is kept because of backward-compatibility so far, and will be deprecated in some future release.

`pgstattuple_approx(regclass)` returns record

`pgstattuple_approx` is a faster alternative to `pgstattuple` that returns approximate results. The argument is the target relation’s name or OID. For example:

```sql
test=> SELECT * FROM pgstattuple_approx('pg_catalog.pg_proc'::regclass);
- [ RECORD 1 ]---------
table_len | 573440
scanned_percent | 2
```
approx_tuple_count | 2740
approx_tuple_len | 561210
approx_tuple_percent | 97.87
dead_tuple_count | 0
dead_tuple_len | 0
dead_tuple_percent | 0
approx_free_space | 11996
approx_free_percent | 2.09

The output columns are described in Table F-24.

Whereas `pgstattuple` always performs a full-table scan and returns an exact count of live and dead tuples (and their sizes) and free space, `pgstattuple_approx` tries to avoid the full-table scan and returns exact dead tuple statistics along with an approximation of the number and size of live tuples and free space.

It does this by skipping pages that have only visible tuples according to the visibility map (if a page has the corresponding VM bit set, then it is assumed to contain no dead tuples). For such pages, it derives the free space value from the free space map, and assumes that the rest of the space on the page is taken up by live tuples.

For pages that cannot be skipped, it scans each tuple, recording its presence and size in the appropriate counters, and adding up the free space on the page. At the end, it estimates the total number of live tuples based on the number of pages and tuples scanned (in the same way that `VACUUM` estimates `pg_class.reltuples`).

<table>
<thead>
<tr>
<th>Column</th>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>table_len</td>
<td>bigint</td>
<td>Physical relation length in bytes (exact)</td>
</tr>
<tr>
<td>scanned_percent</td>
<td>float8</td>
<td>Percentage of table scanned</td>
</tr>
<tr>
<td>approx_tuple_count</td>
<td>bigint</td>
<td>Number of live tuples (estimated)</td>
</tr>
<tr>
<td>approx_tuple_len</td>
<td>bigint</td>
<td>Total length of live tuples in bytes (estimated)</td>
</tr>
<tr>
<td>approx_tuple_percent</td>
<td>float8</td>
<td>Percentage of live tuples</td>
</tr>
<tr>
<td>dead_tuple_count</td>
<td>bigint</td>
<td>Number of dead tuples (exact)</td>
</tr>
<tr>
<td>dead_tuple_len</td>
<td>bigint</td>
<td>Total length of dead tuples in bytes (exact)</td>
</tr>
<tr>
<td>dead_tuple_percent</td>
<td>float8</td>
<td>Percentage of dead tuples</td>
</tr>
<tr>
<td>approx_free_space</td>
<td>bigint</td>
<td>Total free space in bytes (estimated)</td>
</tr>
<tr>
<td>approx_free_percent</td>
<td>float8</td>
<td>Percentage of free space</td>
</tr>
</tbody>
</table>

In the above output, the free space figures may not match the `pgstattuple` output exactly, because the free space map gives us an exact figure, but is not guaranteed to be accurate to the byte.

F.30.2. Authors

Tatsuo Ishii, Satoshi Nagayasu and Abhijit Menon-Sen
The \texttt{pg\_trgm} module provides functions and operators for determining the similarity of alphanumeric text based on trigram matching, as well as index operator classes that support fast searching for similar strings.

\subsubsection*{F.31.1. Trigram (or Trigraph) Concepts}

A trigram is a group of three consecutive characters taken from a string. We can measure the similarity of two strings by counting the number of trigrams they share. This simple idea turns out to be very effective for measuring the similarity of words in many natural languages.

\textbf{Note:} \texttt{pg\_trgm} ignores non-word characters (non-alphanumerics) when extracting trigrams from a string. Each word is considered to have two spaces prefixed and one space suffixed when determining the set of trigrams contained in the string. For example, the set of trigrams in the string "cat" is \texttt{"c", "ca", "cat", and "at"}. The set of trigrams in the string "foo|bar" is \texttt{"f", "fo", "foo", "oo", "b", "ba", "bar", and "ar"}.

\subsubsection*{F.31.2. Functions and Operators}

The functions provided by the \texttt{pg\_trgm} module are shown in Table F-25, the operators in Table F-26.

\begin{table}[h]
\centering
\begin{tabular}{|l|l|l|}
\hline
Function & Returns & Description \\
\hline
\texttt{similarity(text, text)} & real & Returns a number that indicates how similar the two arguments are. The range of the result is zero (indicating that the two strings are completely dissimilar) to one (indicating that the two strings are identical). \\
\texttt{show_trgm(text)} & \\
\texttt{word_similarity(text, text)} & real & Returns an array of all the trigrams in the given string. (In practice this is seldom useful except for debugging.) \\
\hline
\end{tabular}
\caption{\texttt{pg\_trgm} Functions}
\end{table}
### Appendix F: Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>show_limit()</td>
<td>real</td>
<td>Returns the current similarity threshold used by the % operator. This sets the minimum similarity between two words for them to be considered similar enough to be misspellings of each other, for example (deprecated).</td>
</tr>
<tr>
<td>set_limit(real)</td>
<td>real</td>
<td>Sets the current similarity threshold that is used by the % operator. The threshold must be between 0 and 1 (default is 0.3). Returns the same value passed in (deprecated).</td>
</tr>
</tbody>
</table>

Consider the following example:

```sql
# SELECT word_similarity('word', 'two words');
word_similarity
-----------------
0.8
(1 row)
```

In the first string, the set of trigrams is `{"w"," wo","ord","wor","rd "}`. In the second string, the ordered set of trigrams is `{"t"," tw","two","wo "," w"," wo","wor","ord","rds","ds "}`. The most similar extent of an ordered set of trigrams in the second string is `{"w"," wo","wor","ord"}`, and the similarity is 0.8.

This function returns a value that can be approximately understood as the greatest similarity between the first string and any substring of the second string. However, this function does not add padding to the boundaries of the extent. Thus, the number of additional characters present in the second string is not considered, except for the mismatched word boundaries.

### Table F-26. pg_trgm Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>text % text</td>
<td>boolean</td>
<td>Returns true if its arguments have a similarity that is greater than the current similarity threshold set by pg_trgm.similarity_threshold.</td>
</tr>
<tr>
<td>text &lt;# text</td>
<td>boolean</td>
<td>Returns true if the similarity between the trigram set in the first argument and a continuous extent of an ordered trigram set in the second argument is greater than the current word similarity threshold set by pg_trgm.word_similarity_threshold parameter.</td>
</tr>
</tbody>
</table>
### F.31.3. GUC Parameters

**pg_trgm.similarity_threshold** (`real`)  
Sets the current similarity threshold that is used by the `%` operator. The threshold must be between 0 and 1 (default is 0.3).

**pg_trgm.word_similarity_threshold** (`real`)  
Sets the current word similarity threshold that is used by the `<%` and `%>` operators. The threshold must be between 0 and 1 (default is 0.6).

### F.31.4. Index Support

The `pg_trgm` module provides GiST and GIN index operator classes that allow you to create an index over a text column for the purpose of very fast similarity searches. These index types support the above-described similarity operators, and additionally support trigram-based index searches for `LIKE`, `ILIKE`, `~` and `~*` queries. (These indexes do not support equality nor simple comparison operators, so you may need a regular B-tree index too.)

**Example:**

```sql
CREATE TABLE test_trgm (t text);
CREATE INDEX trgm_idx ON test_trgm USING GIST (t gist_trgm_ops);
```

or

```sql
CREATE INDEX trgm_idx ON test_trgm USING GIN (t gin_trgm_ops);
```

At this point, you will have an index on the `t` column that you can use for similarity searching. A typical query is

```sql
SELECT t, similarity(t, 'word') AS sml
FROM test_trgm
WHERE t % 'word'
ORDER BY sml DESC, t;
```
This will return all values in the text column that are sufficiently similar to word, sorted from best match to worst. The index will be used to make this a fast operation even over very large data sets.

A variant of the above query is

```
SELECT t, t <-> 'word' AS dist
    FROM test_trgm
    ORDER BY dist LIMIT 10;
```

This can be implemented quite efficiently by GiST indexes, but not by GIN indexes. It will usually beat the first formulation when only a small number of the closest matches is wanted.

Also you can use an index on the t column for word similarity. For example:

```
SELECT t, word_similarity('word', t) AS sml
    FROM test_trgm
    WHERE 'word' <<% t
    ORDER BY sml DESC, t;
```

This will return all values in the text column for which there is a continuous extent in the corresponding ordered trigram set that is sufficiently similar to the trigram set of word, sorted from best match to worst. The index will be used to make this a fast operation even over very large data sets.

A variant of the above query is

```
SELECT t, 'word' <-> t AS dist
    FROM test_trgm
    ORDER BY dist LIMIT 10;
```

This can be implemented quite efficiently by GiST indexes, but not by GIN indexes.

Beginning in PostgreSQL 9.1, these index types also support index searches for \texttt{LIKE} and \texttt{ILIKE}, for example

```
SELECT * FROM test_trgm WHERE t LIKE '%foo%bar';
```

The index search works by extracting trigrams from the search string and then looking these up in the index. The more trigrams in the search string, the more effective the index search is. Unlike B-tree based searches, the search string need not be left-anchored.

Beginning in PostgreSQL 9.3, these index types also support index searches for regular-expression matches (\texttt{~} and \texttt{~*} operators), for example

```
SELECT * FROM test_trgm WHERE t ~ '(foo|bar)';
```

The index search works by extracting trigrams from the regular expression and then looking these up in the index. The more trigrams that can be extracted from the regular expression, the more effective the index search is. Unlike B-tree based searches, the search string need not be left-anchored.

For both \texttt{LIKE} and regular-expression searches, keep in mind that a pattern with no extractable trigrams will degenerate to a full-index scan.

The choice between GiST and GIN indexing depends on the relative performance characteristics of GiST and GIN, which are discussed elsewhere.
F.31.5. Text Search Integration

Trigram matching is a very useful tool when used in conjunction with a full text index. In particular it can help to recognize misspelled input words that will not be matched directly by the full text search mechanism.

The first step is to generate an auxiliary table containing all the unique words in the documents:

```sql
CREATE TABLE words AS SELECT word FROM
ts_stat('SELECT to_tsvector("simple", bodytext) FROM documents');
```

where documents is a table that has a text field bodytext that we wish to search. The reason for using the simple configuration with the to_tsvector function, instead of using a language-specific configuration, is that we want a list of the original (unstemmed) words.

Next, create a trigram index on the word column:

```sql
CREATE INDEX words_idx ON words USING GIN (word gin_trgm_ops);
```

Now, a SELECT query similar to the previous example can be used to suggest spellings for misspelled words in user search terms. A useful extra test is to require that the selected words are also of similar length to the misspelled word.

**Note:** Since the words table has been generated as a separate, static table, it will need to be periodically regenerated so that it remains reasonably up-to-date with the document collection. Keeping it exactly current is usually unnecessary.

F.31.6. References

GiST Development Site http://www.sai.msu.su/~megera/postgres/gist/
Tsearch2 Development Site http://www.sai.msu.su/~megera/postgres/gist/tsearch/V2/

F.31.7. Authors

Oleg Bartunov <oleg@sai.msu.su>, Moscow, Moscow University, Russia
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Documentation: Christopher Kings-Lynne

This module is sponsored by Delta-Soft Ltd., Moscow, Russia.

F.32. pg_visibility

The pg_visibility module provides a means for examining the visibility map (VM) and page-level visibility information of a table. It also provides functions to check the integrity of a visibility map and to force it to be rebuilt.
Three different bits are used to store information about page-level visibility. The all-visible bit in the visibility map indicates that every tuple in the corresponding page of the relation is visible to every current and future transaction. The all-frozen bit in the visibility map indicates that every tuple in the page is frozen; that is, no future vacuum will need to modify the page until such time as a tuple is inserted, updated, deleted, or locked on that page. The page header’s PD_ALL_VISIBLE bit has the same meaning as the all-visible bit in the visibility map, but is stored within the data page itself rather than in a separate data structure. These two bits will normally agree, but the page’s all-visible bit can sometimes be set while the visibility map bit is clear after a crash recovery. The reported values can also disagree because of a change that occurs after \texttt{pg_visibility} examines the visibility map and before it examines the data page. Any event that causes data corruption can also cause these bits to disagree.

Functions that display information about PD_ALL_VISIBLE bits are much more costly than those that only consult the visibility map, because they must read the relation’s data blocks rather than only the (much smaller) visibility map. Functions that check the relation’s data blocks are similarly expensive.

\section*{F.32.1. Functions}

\begin{verbatim}
pg_visibility_map(relation regclass, blkno bigint, all_visible OUT boolean, all_frozen OUT boolean) returns record

  Returns the all-visible and all-frozen bits in the visibility map for the given block of the given relation.

pg_visibility(relation regclass, blkno bigint, all_visible OUT boolean, all_frozen OUT boolean, pd_all_visible OUT boolean) returns record

  Returns the all-visible and all-frozen bits in the visibility map for the given block of the given relation, plus the PD_ALL_VISIBLE bit of that block.

pg_visibility_map(relation regclass, blkno OUT bigint, all_visible OUT boolean, all_frozen OUT boolean) returns setof record

  Returns the all-visible and all-frozen bits in the visibility map for each block of the given relation.

pg_visibility(relation regclass, blkno OUT bigint, all_visible OUT boolean, all_frozen OUT boolean, pd_all_visible OUT boolean) returns setof record

  Returns the all-visible and all-frozen bits in the visibility map for each block of the given relation, plus the PD_ALL_VISIBLE bit of each block.

pg_visibility_map_summary(relation regclass, all_visible OUT bigint, all_frozen OUT bigint) returns record

  Returns the number of all-visible pages and the number of all-frozen pages in the relation according to the visibility map.

pg_check_frozen(relation regclass, t_ctid OUT tid) returns setof tid

  Returns the TIDs of non-frozen tuples stored in pages marked all-frozen in the visibility map. If this function returns a non-empty set of TIDs, the visibility map is corrupt.

pg_check_visible(relation regclass, t_ctid OUT tid) returns setof tid

  Returns the TIDs of non-all-visible tuples stored in pages marked all-visible in the visibility map. If this function returns a non-empty set of TIDs, the visibility map is corrupt.
\end{verbatim}
Appendix F. Additional Supplied Modules

pg_truncate_visibility_map(relation regclass) returns void

Truncates the visibility map for the given relation. This function is useful if you believe that the visibility map for the relation is corrupt and wish to force rebuilding it. The first VACUUM executed on the given relation after this function is executed will scan every page in the relation and rebuild the visibility map. (Until that is done, queries will treat the visibility map as containing all zeroes.)

By default, these functions are executable only by superusers.

F.32.2. Author

Robert Haas <rhaas@postgresql.org>

F.33. postgres_fdw

The postgres_fdw module provides the foreign-data wrapper postgres_fdw, which can be used to access data stored in external PostgreSQL servers.

The functionality provided by this module overlaps substantially with the functionality of the older dblink module. But postgres_fdw provides more transparent and standards-compliant syntax for accessing remote tables, and can give better performance in many cases.

To prepare for remote access using postgres_fdw:

1. Install the postgres_fdw extension using CREATE EXTENSION.
2. Create a foreign server object, using CREATE SERVER, to represent each remote database you want to connect to. Specify connection information, except user and password, as options of the server object.
3. Create a user mapping, using CREATE USER MAPPING, for each database user you want to allow to access each foreign server. Specify the remote user name and password to use as user and password options of the user mapping.
4. Create a foreign table, using CREATE FOREIGN TABLE or IMPORT FOREIGN SCHEMA, for each remote table you want to access. The columns of the foreign table must match the referenced remote table. You can, however, use table and/or column names different from the remote table’s, if you specify the correct remote names as options of the foreign table object.

Now you need only SELECT from a foreign table to access the data stored in its underlying remote table. You can also modify the remote table using INSERT, UPDATE, or DELETE. (Of course, the remote user you have specified in your user mapping must have privileges to do these things.)

Note that postgres_fdw currently lacks support for INSERT statements with an ON CONFLICT DO UPDATE clause. However, the ON CONFLICT DO NOTHING clause is supported, provided a unique index inference specification is omitted.

It is generally recommended that the columns of a foreign table be declared with exactly the same data types, and collations if applicable, as the referenced columns of the remote table. Although postgres_fdw is currently rather forgiving about performing data type conversions at need, surprising semantic anomalies may arise when types or collations do not match, due to the remote server interpreting WHERE clauses slightly differently from the local server.
Note that a foreign table can be declared with fewer columns, or with a different column order, than its underlying remote table has. Matching of columns to the remote table is by name, not position.

F.33.1. FDW Options of postgres_fdw

F.33.1.1. Connection Options

A foreign server using the `postgres_fdw` foreign data wrapper can have the same options that libpq accepts in connection strings, as described in Section 32.1.2, except that these options are not allowed:

- `user` and `password` (specify these in a user mapping, instead)
- `client_encoding` (this is automatically set from the local server encoding)
- `fallback_application_name` (always set to `postgres_fdw`)

Only superusers may connect to foreign servers without password authentication, so always specify the `password` option for user mappings belonging to non-superusers.

F.33.1.2. Object Name Options

These options can be used to control the names used in SQL statements sent to the remote PostgreSQL server. These options are needed when a foreign table is created with names different from the underlying remote table’s names.

`schema_name`

This option, which can be specified for a foreign table, gives the schema name to use for the foreign table on the remote server. If this option is omitted, the name of the foreign table’s schema is used.

`table_name`

This option, which can be specified for a foreign table, gives the table name to use for the foreign table on the remote server. If this option is omitted, the foreign table’s name is used.

`column_name`

This option, which can be specified for a column of a foreign table, gives the column name to use for the column on the remote server. If this option is omitted, the column’s name is used.

F.33.1.3. Cost Estimation Options

`postgres_fdw` retrieves remote data by executing queries against remote servers, so ideally the estimated cost of scanning a foreign table should be whatever it costs to be done on the remote server, plus some overhead for communication. The most reliable way to get such an estimate is to ask the remote server and then add something for overhead — but for simple queries, it may not be worth the cost of an additional remote query to get a cost estimate. So `postgres_fdw` provides the following options to control how cost estimation is done:
Appendix F. Additional Supplied Modules

use_remote_estimate

This option, which can be specified for a foreign table or a foreign server, controls whether postgres_fdw issues remote EXPLAIN commands to obtain cost estimates. A setting for a foreign table overrides any setting for its server, but only for that table. The default is false.

fdw_startup_cost

This option, which can be specified for a foreign server, is a numeric value that is added to the estimated startup cost of any foreign-table scan on that server. This represents the additional overhead of establishing a connection, parsing and planning the query on the remote side, etc. The default value is 100.

fdw_tuple_cost

This option, which can be specified for a foreign server, is a numeric value that is used as extra cost per-tuple for foreign-table scans on that server. This represents the additional overhead of data transfer between servers. You might increase or decrease this number to reflect higher or lower network delay to the remote server. The default value is 0.01.

When use_remote_estimate is true, postgres_fdw obtains row count and cost estimates from the remote server and then adds fdw_startup_cost and fdw_tuple_cost to the cost estimates. When use_remote_estimate is false, postgres_fdw performs local row count and cost estimation and then adds fdw_startup_cost and fdw_tuple_cost to the cost estimates. This local estimation is unlikely to be very accurate unless local copies of the remote table’s statistics are available. Running ANALYZE on the foreign table is the way to update the local statistics; this will perform a scan of the remote table and then calculate and store statistics just as though the table were local. Keeping local statistics can be a useful way to reduce per-query planning overhead for a remote table — but if the remote table is frequently updated, the local statistics will soon be obsolete.

F.33.1.4. Remote Execution Options

By default, only WHERE clauses using built-in operators and functions will be considered for execution on the remote server. Clauses involving non-built-in functions are checked locally after rows are fetched. If such functions are available on the remote server and can be relied on to produce the same results as they do locally, performance can be improved by sending such WHERE clauses for remote execution. This behavior can be controlled using the following option:

extensions

This option is a comma-separated list of names of PostgreSQL extensions that are installed, in compatible versions, on both the local and remote servers. Functions and operators that are immutable and belong to a listed extension will be considered shippable to the remote server. This option can only be specified for foreign servers, not per-table.

When using the extensions option, it is the user’s responsibility that the listed extensions exist and behave identically on both the local and remote servers. Otherwise, remote queries may fail or behave unexpectedly.

fetch_size

This option specifies the number of rows postgres_fdw should get in each fetch operation. It can be specified for a foreign table or a foreign server. The option specified on a table overrides an option specified for the server. The default is 100.
F.33.1.5. Updatability Options

By default all foreign tables using `postgres_fdw` are assumed to be updatable. This may be overridden using the following option:

```
updatable
```

This option controls whether `postgres_fdw` allows foreign tables to be modified using `INSERT`, `UPDATE` and `DELETE` commands. It can be specified for a foreign table or a foreign server. A table-level option overrides a server-level option. The default is `true`.

Of course, if the remote table is not in fact updatable, an error would occur anyway. Use of this option primarily allows the error to be thrown locally without querying the remote server. Note however that the `information_schema` views will report a `postgres_fdw` foreign table to be updatable (or not) according to the setting of this option, without any check of the remote server.

F.33.1.6. Importing Options

`postgres_fdw` is able to import foreign table definitions using `IMPORT FOREIGN SCHEMA`. This command creates foreign table definitions on the local server that match tables or views present on the remote server. If the remote tables to be imported have columns of user-defined data types, the local server must have compatible types of the same names.

Importing behavior can be customized with the following options (given in the `IMPORT FOREIGN SCHEMA` command):

```
import_collate
```

This option controls whether column `COLLATE` options are included in the definitions of foreign tables imported from a foreign server. The default is `true`. You might need to turn this off if the remote server has a different set of collation names than the local server does, which is likely to be the case if it’s running on a different operating system.

```
import_default
```

This option controls whether column `DEFAULT` expressions are included in the definitions of foreign tables imported from a foreign server. The default is `false`. If you enable this option, be wary of defaults that might get computed differently on the local server than they would be on the remote server; `nextval()` is a common source of problems. The `IMPORT` will fail altogether if an imported default expression uses a function or operator that does not exist locally.

```
import_not_null
```

This option controls whether column `NOT NULL` constraints are included in the definitions of foreign tables imported from a foreign server. The default is `true`.

Note that constraints other than `NOT NULL` will never be imported from the remote tables. Although PostgreSQL does support `CHECK` constraints on foreign tables, there is no provision for importing them automatically, because of the risk that a constraint expression could evaluate differently on the local and remote servers. Any such inconsistency in the behavior of a `CHECK` constraint could lead to hard-to-detect errors in query optimization. So if you wish to import `CHECK` constraints, you must do so manually, and you should verify the semantics of each one carefully. For more detail about the treatment of `CHECK` constraints on foreign tables, see `CREATE FOREIGN TABLE`.
F.33.2. Connection Management

`postgres_fdw` establishes a connection to a foreign server during the first query that uses a foreign table associated with the foreign server. This connection is kept and re-used for subsequent queries in the same session. However, if multiple user identities (user mappings) are used to access the foreign server, a connection is established for each user mapping.

F.33.3. Transaction Management

During a query that references any remote tables on a foreign server, `postgres_fdw` opens a transaction on the remote server if one is not already open corresponding to the current local transaction. The remote transaction is committed or aborted when the local transaction commits or aborts. Savepoints are similarly managed by creating corresponding remote savepoints.

The remote transaction uses `SERIALIZABLE` isolation level when the local transaction has `SERIALIZABLE` isolation level; otherwise it uses `REPEATABLE READ` isolation level. This choice ensures that if a query performs multiple table scans on the remote server, it will get snapshot-consistent results for all the scans. A consequence is that successive queries within a single transaction will see the same data from the remote server, even if concurrent updates are occurring on the remote server due to other activities. That behavior would be expected anyway if the local transaction uses `SERIALIZABLE` or `REPEATABLE READ` isolation level, but it might be surprising for a `READ COMMITTED` local transaction. A future PostgreSQL release might modify these rules.

Note that it is currently not supported by `postgres_fdw` to prepare the remote transaction for two-phase commit.

F.33.4. Remote Query Optimization

`postgres_fdw` attempts to optimize remote queries to reduce the amount of data transferred from foreign servers. This is done by sending query `WHERE` clauses to the remote server for execution, and by not retrieving table columns that are not needed for the current query. To reduce the risk of misexecution of queries, `WHERE` clauses are not sent to the remote server unless they use only data types, operators, and functions that are built-in or belong to an extension that’s listed in the foreign server’s `extensions` option. Operators and functions in such clauses must be `IMMUTABLE` as well. For an `UPDATE` or `DELETE` query, `postgres_fdw` attempts to optimize the query execution by sending the whole query to the remote server if there are no query `WHERE` clauses that cannot be sent to the remote server, no local joins for the query, no row-level local `BEFORE` or `AFTER` triggers on the target table, and no `CHECK OPTION` constraints from parent views. In `UPDATE`, expressions to assign to target columns must use only built-in data types, `IMMUTABLE` operators, or `IMMUTABLE` functions, to reduce the risk of misexecution of the query.

When `postgres_fdw` encounters a join between foreign tables on the same foreign server, it sends the entire join to the foreign server, unless for some reason it believes that it will be more efficient to fetch rows from each table individually, or unless the table references involved are subject to different user mappings. While sending the `JOIN` clauses, it takes the same precautions as mentioned above for the `WHERE` clauses.

The query that is actually sent to the remote server for execution can be examined using `EXPLAIN VERBOSE`.

2472
F.33.5. Remote Query Execution Environment

In the remote sessions opened by `postgres_fdw`, the search_path parameter is set to just `pg_catalog`, so that only built-in objects are visible without schema qualification. This is not an issue for queries generated by `postgres_fdw` itself, because it always supplies such qualification. However, this can pose a hazard for functions that are executed on the remote server via triggers or rules on remote tables. For example, if a remote table is actually a view, any functions used in that view will be executed with the restricted search path. It is recommended to schema-qualify all names in such functions, or else attach SET search_path options (see CREATE FUNCTION) to such functions to establish their expected search path environment.

`postgres_fdw` likewise establishes remote session settings for the parameters TimeZone, DateStyle, IntervalStyle, and extra_float_digits. These are less likely to be problematic than search_path, but can be handled with function SET options if the need arises.

It is not recommended that you override this behavior by changing the session-level settings of these parameters; that is likely to cause `postgres_fdw` to malfunction.

F.33.6. Cross-Version Compatibility

`postgres_fdw` can be used with remote servers dating back to PostgreSQL 8.3. Read-only capability is available back to 8.1. A limitation however is that `postgres_fdw` generally assumes that immutable built-in functions and operators are safe to send to the remote server for execution, if they appear in a WHERE clause for a foreign table. Thus, a built-in function that was added since the remote server’s release might be sent to it for execution, resulting in “function does not exist” or a similar error. This type of failure can be worked around by rewriting the query, for example by embedding the foreign table reference in a sub-SELECT with OFFSET 0 as an optimization fence, and placing the problematic function or operator outside the sub-SELECT.

F.33.7. Examples

Here is an example of creating a foreign table with `postgres_fdw`. First install the extension:

```
CREATE EXTENSION postgres_fdw;
```

Then create a foreign server using CREATE SERVER. In this example we wish to connect to a PostgreSQL server on host `192.83.123.89` listening on port `5432`. The database to which the connection is made is named `foreign_db` on the remote server:

```
CREATE SERVER foreign_server
FOREIGN DATA WRAPPER postgres_fdw
OPTIONS (host '192.83.123.89', port '5432', dbname 'foreign_db');
```

A user mapping, defined with CREATE USER MAPPING, is needed as well to identify the role that will be used on the remote server:

```
CREATE USER MAPPING FOR local_user
SERVER foreign_server
OPTIONS (user 'foreign_user', password 'password');
```
Now it is possible to create a foreign table with CREATE FOREIGN TABLE. In this example we wish to access the table named `some_schema.some_table` on the remote server. The local name for it will be `foreign_table`:

```
CREATE FOREIGN TABLE foreign_table (
    id integer NOT NULL,
    data text
) SERVER foreign_server
OPTIONS (schema_name 'some_schema', table_name 'some_table');
```

It’s essential that the data types and other properties of the columns declared in CREATE FOREIGN TABLE match the actual remote table. Column names must match as well, unless you attach column_name options to the individual columns to show how they are named in the remote table. In many cases, use of IMPORT FOREIGN SCHEMA is preferable to constructing foreign table definitions manually.

F.33.8. Author

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F.34. seg

This module implements a data type `seg` for representing line segments, or floating point intervals. `seg` can represent uncertainty in the interval endpoints, making it especially useful for representing laboratory measurements.

F.34.1. Rationale

The geometry of measurements is usually more complex than that of a point in a numeric continuum. A measurement is usually a segment of that continuum with somewhat fuzzy limits. The measurements come out as intervals because of uncertainty and randomness, as well as because the value being measured may naturally be an interval indicating some condition, such as the temperature range of stability of a protein.

Using just common sense, it appears more convenient to store such data as intervals, rather than pairs of numbers. In practice, it even turns out more efficient in most applications.

Further along the line of common sense, the fuzziness of the limits suggests that the use of traditional numeric data types leads to a certain loss of information. Consider this: your instrument reads 6.50, and you input this reading into the database. What do you get when you fetch it? Watch:

```
test=> select 6.50 :: float8 as "pH";
       pH
  6.5
   (1 row)
```

In the world of measurements, 6.50 is not the same as 6.5. It may sometimes be critically different. The experimenters usually write down (and publish) the digits they trust. 6.50 is actually a fuzzy interval contained within a bigger and even fuzzier interval, 6.5, with their center points being (probably) the
only common feature they share. We definitely do not want such different data items to appear the same.

Conclusion? It is nice to have a special data type that can record the limits of an interval with arbitrarily variable precision. Variable in the sense that each data element records its own precision.

Check this out:

test=> select '6.25 .. 6.50'::seg as "pH";
\n  pH
--------
  6.25 .. 6.50
(1 row)

F.34.2. Syntax

The external representation of an interval is formed using one or two floating-point numbers joined by the range operator (.. or ...). Alternatively, it can be specified as a center point plus or minus a deviation. Optional certainty indicators (<, > or ~) can be stored as well. (Certainty indicators are ignored by all the built-in operators, however.) Table F-27 gives an overview of allowed representations; Table F-28 shows some examples.

In Table F-27, x, y, and delta denote floating-point numbers. x and y, but not delta, can be preceded by a certainty indicator.

Table F-27. seg External Representations

<table>
<thead>
<tr>
<th>Representation</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>x</td>
<td>Single value (zero-length interval)</td>
</tr>
<tr>
<td>x .. y</td>
<td>Interval from x to y</td>
</tr>
<tr>
<td>x (+-) delta</td>
<td>Interval from x - delta to x + delta</td>
</tr>
<tr>
<td>x ..</td>
<td>Open interval with lower bound x</td>
</tr>
<tr>
<td>.. x</td>
<td>Open interval with upper bound x</td>
</tr>
</tbody>
</table>

Table F-28. Examples of Valid seg Input

<table>
<thead>
<tr>
<th>Input</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>5.0</td>
<td>Creates a zero-length segment (a point, if you will)</td>
</tr>
<tr>
<td>~5.0</td>
<td>Creates a zero-length segment and records ~ in the data. ~ is ignored by seg operations, but is preserved as a comment.</td>
</tr>
<tr>
<td>&lt;5.0</td>
<td>Creates a point at 5.0. &lt; is ignored but is preserved as a comment.</td>
</tr>
<tr>
<td>&gt;5.0</td>
<td>Creates a point at 5.0. &gt; is ignored but is preserved as a comment.</td>
</tr>
<tr>
<td>5(+-)0.3</td>
<td>Creates an interval 4.7 .. 5.3. Note that the (+-) notation isn’t preserved.</td>
</tr>
<tr>
<td>50 ..</td>
<td>Everything that is greater than or equal to 50</td>
</tr>
<tr>
<td>.. 0</td>
<td>Everything that is less than or equal to 0</td>
</tr>
</tbody>
</table>
1.5e-2 .. 2E-2 | Creates an interval 0.015 .. 0.02
1 ... 2 | The same as 1..2, or 1 .. 2, or 1..2 (spaces around the range operator are ignored)

Because ... is widely used in data sources, it is allowed as an alternative spelling of ... Unfortunately, this creates a parsing ambiguity: it is not clear whether the upper bound in 0...23 is meant to be 23 or 0.23. This is resolved by requiring at least one digit before the decimal point in all numbers in seg input.

As a sanity check, seg rejects intervals with the lower bound greater than the upper, for example 5 .. 2.

F.34.3. Precision

seg values are stored internally as pairs of 32-bit floating point numbers. This means that numbers with more than 7 significant digits will be truncated.

Numbers with 7 or fewer significant digits retain their original precision. That is, if your query returns 0.00, you will be sure that the trailing zeroes are not the artifacts of formatting: they reflect the precision of the original data. The number of leading zeroes does not affect precision: the value 0.0067 is considered to have just 2 significant digits.

F.34.4. Usage

The seg module includes a GiST index operator class for seg values. The operators supported by the GiST operator class are shown in Table F-29.

Table F-29. Seg GiST Operators

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[a, b] &lt;&lt;&lt; [c, d]</td>
<td>[a, b] is entirely to the left of [c, d]. That is, [a, b] &lt;&lt;&lt; [c, d] is true if b &lt; c and false otherwise.</td>
</tr>
<tr>
<td>[a, b] &gt;&gt;&gt; [c, d]</td>
<td>[a, b] is entirely to the right of [c, d]. That is, [a, b] &gt;&gt;&gt; [c, d] is true if a &gt; d and false otherwise.</td>
</tr>
<tr>
<td>[a, b] &amp;&lt; [c, d]</td>
<td>Overlaps or is left of — This might be better read as “does not extend to right of”. It is true when b &lt;= d.</td>
</tr>
<tr>
<td>[a, b] &amp;&gt; [c, d]</td>
<td>Overlaps or is right of — This might be better read as “does not extend to left of”. It is true when a &gt;= c.</td>
</tr>
<tr>
<td>[a, b] = [c, d]</td>
<td>Same as — The segments [a, b] and [c, d] are identical, that is, a = c and b = d.</td>
</tr>
<tr>
<td>[a, b] &amp;&amp; [c, d]</td>
<td>The segments [a, b] and [c, d] overlap.</td>
</tr>
<tr>
<td>[a, b] @&gt; [c, d]</td>
<td>The segment [a, b] contains the segment [c, d], that is, a &lt;= c and b &gt;= d.</td>
</tr>
<tr>
<td>[a, b] &lt;@ [c, d]</td>
<td>The segment [a, b] is contained in [c, d], that is, a &gt;= c and b &lt;= d.</td>
</tr>
</tbody>
</table>
(Before PostgreSQL 8.2, the containment operators @> and <@ were respectively called @ and ~. These names are still available, but are deprecated and will eventually be retired. Notice that the old names are reversed from the convention formerly followed by the core geometric data types!)

The standard B-tree operators are also provided, for example

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>[a, b] &lt; [c, d]</td>
<td>Less than</td>
</tr>
<tr>
<td>[a, b] &gt; [c, d]</td>
<td>Greater than</td>
</tr>
</tbody>
</table>

These operators do not make a lot of sense for any practical purpose but sorting. These operators first compare (a) to (c), and if these are equal, compare (b) to (d). That results in reasonably good sorting in most cases, which is useful if you want to use ORDER BY with this type.

**F.34.5. Notes**

For examples of usage, see the regression test sql/seg.sql.

The mechanism that converts (+-) to regular ranges isn’t completely accurate in determining the number of significant digits for the boundaries. For example, it adds an extra digit to the lower boundary if the resulting interval includes a power of ten:

```
postgres=> select ‘10(+-)1’:seg as seg;
  seg
---------
 9.0 .. 11 -- should be: 9 .. 11
```

The performance of an R-tree index can largely depend on the initial order of input values. It may be very helpful to sort the input table on the seg column; see the script sort-segments.pl for an example.

**F.34.6. Credits**

Original author: Gene Selkov, Jr. <selkovjr@mcs.anl.gov>, Mathematics and Computer Science Division, Argonne National Laboratory.

My thanks are primarily to Prof. Joe Hellerstein (http://db.cs.berkeley.edu/jmh/) for elucidating the gist of the GiST (http://gist.cs.berkeley.edu/). I am also grateful to all Postgres developers, present and past, for enabling myself to create my own world and live undisturbed in it. And I would like to acknowledge my gratitude to Argonne Lab and to the U.S. Department of Energy for the years of faithful support of my database research.

**F.35. sepgsql**

sepgsql is a loadable module that supports label-based mandatory access control (MAC) based on SELinux security policy.
Warning
The current implementation has significant limitations, and does not enforce mandatory access control for all actions. See Section F.35.7.

F.35.1. Overview
This module integrates with SELinux to provide an additional layer of security checking above and beyond what is normally provided by PostgreSQL. From the perspective of SELinux, this module allows PostgreSQL to function as a user-space object manager. Each table or function access initiated by a DML query will be checked against the system security policy. This check is in addition to the usual SQL permissions checking performed by PostgreSQL.

SELinux access control decisions are made using security labels, which are represented by strings such as `system_u:object_r:sepgsql_table_t:s0`. Each access control decision involves two labels: the label of the subject attempting to perform the action, and the label of the object on which the operation is to be performed. Since these labels can be applied to any sort of object, access control decisions for objects stored within the database can be (and, with this module, are) subjected to the same general criteria used for objects of any other type, such as files. This design is intended to allow a centralized security policy to protect information assets independent of the particulars of how those assets are stored.

The SECURITY LABEL statement allows assignment of a security label to a database object.

F.35.2. Installation
`sepgsql` can only be used on Linux 2.6.28 or higher with SELinux enabled. It is not available on any other platform. You will also need libselinux 2.1.10 or higher and selinux-policy 3.9.13 or higher (although some distributions may backport the necessary rules into older policy versions).

The `sestatus` command allows you to check the status of SELinux. A typical display is:

```
$ sestatus
SELinux status: enabled
SELinuxfs mount: /selinux
Current mode: enforcing
Mode from config file: enforcing
Policy version: 24
Policy from config file: targeted
```

If SELinux is disabled or not installed, you must set that product up first before installing this module.

To build this module, include the option `--with-selinux` in your PostgreSQL `configure` command. Be sure that the `libselinux-devel` RPM is installed at build time.

To use this module, you must include `sepgsql` in the `shared_preload_libraries` parameter in `postgresql.conf`. The module will not function correctly if loaded in any other manner. Once the module is loaded, you should execute `sepgsql.sql` in each database. This will install functions needed for security label management, and assign initial security labels.

Here is an example showing how to initialize a fresh database cluster with `sepgsql` functions and security labels installed. Adjust the paths shown as appropriate for your installation:

```
$ export PGDATA=/path/to/data/directory
$ initdb
```
Appendix F. Additional Supplied Modules

$ vi $PGDATA/postgresql.conf
change
  #shared_preload_libraries = " # (change requires restart)
to
  shared_preload_libraries = 'sepgsql' # (change requires restart)

$ for DBNAME in template0 template1 postgres; do
  postgres --single -F -c exit_on_error=true $DBNAME \<
    /usr/local/pgsql/share/contrib/sepgsql.sql >/dev/null
done

Please note that you may see some or all of the following notifications depending on the particular versions you have of libselinux and selinux-policy:

/etc/selinux/targeted/contexts/sepgsql_contexts: line 33 has invalid object type db_blobs
/etc/selinux/targeted/contexts/sepgsql_contexts: line 36 has invalid object type db_language
/etc/selinux/targeted/contexts/sepgsql_contexts: line 37 has invalid object type db_language
/etc/selinux/targeted/contexts/sepgsql_contexts: line 38 has invalid object type db_language
/etc/selinux/targeted/contexts/sepgsql_contexts: line 39 has invalid object type db_language
/etc/selinux/targeted/contexts/sepgsql_contexts: line 40 has invalid object type db_language

These messages are harmless and should be ignored.

If the installation process completes without error, you can now start the server normally.

F.35.3. Regression Tests

Due to the nature of SELinux, running the regression tests for sepgsql requires several extra configuration steps, some of which must be done as root. The regression tests will not be run by an ordinary make check or make installcheck command; you must set up the configuration and then invoke the test script manually. The tests must be run in the contrib/sepgsql directory of a configured PostgreSQL build tree. Although they require a build tree, the tests are designed to be executed against an installed server, that is they are comparable to make installcheck not make check.

First, set up sepgsql in a working database according to the instructions in Section F.35.2. Note that the current operating system user must be able to connect to the database as superuser without password authentication.

Second, build and install the policy package for the regression test. The sepgsql-regtest policy is a special purpose policy package which provides a set of rules to be allowed during the regression tests. It should be built from the policy source file sepgsql-regtest.te, which is done using make with a Makefile supplied by SELinux. You will need to locate the appropriate Makefile on your system; the path shown below is only an example. Once built, install this policy package using the semodule command, which loads supplied policy packages into the kernel. If the package is correctly installed, semodule -l should list sepgsql-regtest as an available policy package:

$ cd .../contrib/sepgsql
$ make -f /usr/share/selinux/devel/Makefile
$ sudo semodule -u sepgsql-regtest.pp
$ sudo semodule -l | grep sepgsql
sepgsql-regtest 1.07

Third, turn on sepgsql_regression_test_mode. For security reasons, the rules in sepgsql-regtest are not enabled by default; the sepgsql_regression_test_mode parameter enables the rules needed to launch the regression tests. It can be turned on using the setsebool command:

$ setsebool -P sepgsql_regression_test_mode=on
Appendix F. Additional Supplied Modules

$ sudo setsebool sepgsql_regression_test_mode on
$ getsebool sepgsql_regression_test_mode
sepgsql_regression_test_mode --> on

Fourth, verify your shell is operating in the unconfined_t domain:

$ id -Z
unconfined_u:unconfined_r:unconfined_t:s0-s0:c0.c1023

See Section F.35.8 for details on adjusting your working domain, if necessary.

Finally, run the regression test script:

$ ./test_sepgsql

This script will attempt to verify that you have done all the configuration steps correctly, and then it will run the regression tests for the sepgsql module.

After completing the tests, it’s recommended you disable the sepgsql_regression_test_mode parameter:

$ sudo setsebool sepgsql_regression_test_mode off

You might prefer to remove the sepgsql-regtest policy entirely:

$ sudo semodule -r sepgsql-regtest

F.35.4. GUC Parameters

sepgsql.permissive (boolean)

This parameter enables sepgsql to function in permissive mode, regardless of the system setting. The default is off. This parameter can only be set in the postgresql.conf file or on the server command line.

When this parameter is on, sepgsql functions in permissive mode, even if SELinux in general is working in enforcing mode. This parameter is primarily useful for testing purposes.

sepgsql.debug_audit (boolean)

This parameter enables the printing of audit messages regardless of the system policy settings. The default is off, which means that messages will be printed according to the system settings.

The security policy of SELinux also has rules to control whether or not particular accesses are logged. By default, access violations are logged, but allowed accesses are not.

This parameter forces all possible logging to be turned on, regardless of the system policy.

F.35.5. Features

F.35.5.1. Controlled Object Classes

The security model of SELinux describes all the access control rules as relationships between a subject entity (typically, a client of the database) and an object entity (such as a database object), each of which is identified by a security label. If access to an unlabeled object is attempted, the object is treated as if it were assigned the label unlabeled_t.
Currently, sepgsql allows security labels to be assigned to schemas, tables, columns, sequences, views, and functions. When sepgsql is in use, security labels are automatically assigned to supported database objects at creation time. This label is called a default security label, and is decided according to the system security policy, which takes as input the creator’s label, the label assigned to the new object’s parent object and optionally name of the constructed object.

A new database object basically inherits the security label of the parent object, except when the security policy has special rules known as type-transition rules, in which case a different label may be applied. For schemas, the parent object is the current database; for tables, sequences, views, and functions, it is the containing schema; for columns, it is the containing table.

**F.35.5.2. DML Permissions**

For tables, `db_table:select`, `db_table:insert`, `db_table:update` or `db_table:delete` are checked for all the referenced target tables depending on the kind of statement; in addition, `db_table:select` is also checked for all the tables that contain columns referenced in the WHERE or RETURNING clause, as a data source for UPDATE, and so on.

Column-level permissions will also be checked for each referenced column. `db_column:select` is checked on not only the columns being read using SELECT, but those being referenced in other DML statements; `db_column:update` or `db_column:insert` will also be checked for columns being modified by UPDATE or INSERT.

For example, consider:

```
UPDATE t1 SET x = 2, y = md5sum(y) WHERE z = 100;
```

Here, `db_column:update` will be checked for `t1.x`, since it is being updated, `db_column:{select update}` will be checked for `t1.y`, since it is both updated and referenced, and `db_column:select` will be checked for `t1.z`, since it is only referenced. `db_table:{select update}` will also be checked at the table level.

For sequences, `db_sequence:get_value` is checked when we reference a sequence object using SELECT; however, note that we do not currently check permissions on execution of corresponding functions such as lastval().

For views, `db_view:expand` will be checked, then any other required permissions will be checked on the objects being expanded from the view, individually.

For functions, `db_procedure:{execute}` will be checked when user tries to execute a function as a part of query, or using fast-path invocation. If this function is a trusted procedure, it also checks `db_procedure:{entrypoint}` permission to check whether it can perform as entry point of trusted procedure.

In order to access any schema object, `db_schema:search` permission is required on the containing schema. When an object is referenced without schema qualification, schemas on which this permission is not present will not be searched (just as if the user did not have USAGE privilege on the schema). If an explicit schema qualification is present, an error will occur if the user does not have the requisite permission on the named schema.

The client must be allowed to access all referenced tables and columns, even if they originated from views which were then expanded, so that we apply consistent access control rules independent of the manner in which the table contents are referenced.

The default database privilege system allows database superusers to modify system catalogs using DML commands, and reference or modify toast tables. These operations are prohibited when sepgsql is enabled.
Appendix F. Additional Supplied Modules

F.35.5.3. DDL Permissions

SELinux defines several permissions to control common operations for each object type; such as creation, alter, drop and relabel of security label. In addition, several object types have special permissions to control their characteristic operations; such as addition or deletion of name entries within a particular schema.

Creating a new database object requires create permission. SE Linux will grant or deny this permission based on the client’s security label and the proposed security label for the new object. In some cases, additional privileges are required:

- CREATE DATABASE additionally requires getattr permission for the source or template database.
- Creating a schema object additionally requires add_name permission on the parent schema.
- Creating a table additionally requires permission to create each individual table column, just as if each table column were a separate top-level object.
- Creating a function marked as LEAKPROOF additionally requires install permission. (This permission is also checked when LEAKPROOF is set for an existing function.)

When DROP command is executed, drop will be checked on the object being removed. Permissions will be also checked for objects dropped indirectly via CASCADE. Deletion of objects contained within a particular schema (tables, views, sequences and procedures) additionally requires remove_name on the schema.

When ALTER command is executed, setattr will be checked on the object being modified for each object types, except for subsidiary objects such as the indexes or triggers of a table, where permissions are instead checked on the parent object. In some cases, additional permissions are required:

- Moving an object to a new schema additionally requires remove_name permission on the old schema and add_name permission on the new one.
- Setting the LEAKPROOF attribute on a function requires install permission.
- Using SECURITY LABEL on an object additionally requires relabelfrom permission for the object in conjunction with its old security label and relabelto permission for the object in conjunction with its new security label. (In cases where multiple label providers are installed and the user tries to set a security label, but it is not managed by SE Linux, only setattr should be checked here. This is currently not done due to implementation restrictions.)

F.35.5.4. Trusted Procedures

Trusted procedures are similar to security definer functions or setuid commands. SE Linux provides a feature to allow trusted code to run using a security label different from that of the client, generally for the purpose of providing highly controlled access to sensitive data (e.g. rows might be omitted, or the precision of stored values might be reduced). Whether or not a function acts as a trusted procedure is controlled by its security label and the operating system security policy. For example:

```
postgres=# CREATE TABLE customer {
    cid    int primary key,
    cname  text,
    credit text
};
CREATE TABLE
postgres=# SECURITY LABEL ON COLUMN customer.credit
```

2482
Appendix F. Additional Supplied Modules

IS 'system_u:object_r:sepgsql_secret_table_t:s0';

postgres=# CREATE FUNCTION show_credit(int) RETURNS text
AS 'SELECT regexp_replace(credit, "-[0-9]+$", "-xxxx", "g")
FROM customer WHERE cid = $1'
LANGUAGE sql;

CREATE FUNCTION
postgres=# SECURITY LABEL ON FUNCTION show_credit(int)
IS 'system_u:object_r:sepgsql_trusted_proc_exec_t:s0';
SECURITY LABEL

The above operations should be performed by an administrative user.

postgres=# SELECT * FROM customer;
ERROR: SELinux: security policy violation

postgres=# SELECT cid, cname, show_credit(cid) FROM customer;

<table>
<thead>
<tr>
<th>cid</th>
<th>cname</th>
<th>show_credit</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>taro</td>
<td>1111-2222-3333-xxxx</td>
</tr>
<tr>
<td>2</td>
<td>hanako</td>
<td>5555-6666-7777-xxxx</td>
</tr>
</tbody>
</table>
(2 rows)

In this case, a regular user cannot reference customer.credit directly, but a trusted procedure show_credit allows the user to print the credit card numbers of customers with some of the digits masked out.

F.35.5.5. Dynamic Domain Transitions

It is possible to use SELinux’s dynamic domain transition feature to switch the security label of the client process, the client domain, to a new context, if that is allowed by the security policy. The client domain needs the setcurrent permission and also dyntransition from the old to the new domain.

Dynamic domain transitions should be considered carefully, because they allow users to switch their label, and therefore their privileges, at their option, rather than (as in the case of a trusted procedure) as mandated by the system. Thus, the dyntransition permission is only considered safe when used to switch to a domain with a smaller set of privileges than the original one. For example:

regression=# select sepgsql_getcon();

sepgsql_getcon
unconfined_u:unconfined_r:unconfined_t:s0-s0:c0.c1023
(1 row)

regression=# SELECT sepgsql_setcon('unconfined_u:unconfined_r:unconfined_t:s0-s0:c1.c4')

sepgsql_setcon
--------------
t
(1 row)

regression=# SELECT sepgsql_setcon('unconfined_u:unconfined_r:unconfined_t:s0-s0:c1.c4')

ERROR: SELinux: security policy violation

In this example above we were allowed to switch from the larger MCS range c1.c1023 to the smaller range c1.c4, but switching back was denied.
Appendix F. Additional Supplied Modules

A combination of dynamic domain transition and trusted procedure enables an interesting use case that fits the typical process life-cycle of connection pooling software. Even if your connection pooling software is not allowed to run most of SQL commands, you can allow it to switch the security label of the client using the `sepgsql_setcon()` function from within a trusted procedure; that should take some credential to authorize the request to switch the client label. After that, this session will have the privileges of the target user, rather than the connection pooler. The connection pooler can later revert the security label change by again using `sepgsql_setcon()` with NULL argument, again invoked from within a trusted procedure with appropriate permissions checks. The point here is that only the trusted procedure actually has permission to change the effective security label, and only does so when given proper credentials. Of course, for secure operation, the credential store (table, procedure definition, or whatever) must be protected from unauthorized access.

F.35.5.6. Miscellaneous

We reject the LOAD command across the board, because any module loaded could easily circumvent security policy enforcement.

F.35.6. Sepgresql Functions

Table F.30 shows the available functions.

Table F.30. Sepgresql Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>sepgsql_getcon()</code></td>
<td>text</td>
<td>Returns the client domain, the current security label of the client.</td>
</tr>
<tr>
<td><code>sepgsql_setcon(text)</code></td>
<td>bool</td>
<td>Switches the client domain of the current session to the new domain, if allowed by the security policy. It also accepts NULL input as a request to transition to the client’s original domain.</td>
</tr>
<tr>
<td><code>sepgsql_mcstrans_in(text)</code></td>
<td>text</td>
<td>Translates the given qualified MLS/MCS range into raw format if the mcstrans daemon is running.</td>
</tr>
<tr>
<td><code>sepgsql_mcstrans_out(text)</code></td>
<td>text</td>
<td>Translates the given raw MLS/MCS range into qualified format if the mcstrans daemon is running.</td>
</tr>
<tr>
<td><code>sepgsql_restorecon(text)</code></td>
<td>bool</td>
<td>Sets up initial security labels for all objects within the current database. The argument may be NULL, or the name of a specfile to be used as alternative of the system default.</td>
</tr>
</tbody>
</table>

F.35.7. Limitations

Data Definition Language (DDL) Permissions

Due to implementation restrictions, some DDL operations do not check permissions.
Appendix F. Additional Supplied Modules

Data Control Language (DCL) Permissions

Due to implementation restrictions, DCL operations do not check permissions.

Row-level access control

PostgreSQL supports row-level access, but sepgsql does not.

Covert channels

sepgsql does not try to hide the existence of a certain object, even if the user is not allowed to reference it. For example, we can infer the existence of an invisible object as a result of primary key conflicts, foreign key violations, and so on, even if we cannot obtain the contents of the object. The existence of a top secret table cannot be hidden; we only hope to conceal its contents.

F.35.8. External Resources

SE-PostgreSQL Introduction

This wiki page provides a brief overview, security design, architecture, administration and upcoming features.

Fedora SELinux User Guide

This document provides a wide spectrum of knowledge to administer SELinux on your systems. It focuses primarily on Fedora, but is not limited to Fedora.

Fedora SELinux FAQ

This document answers frequently asked questions about SELinux. It focuses primarily on Fedora, but is not limited to Fedora.

F.35.9. Author

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F.36. spi

The spi module provides several workable examples of using the Server Programming Interface (SPI) and triggers. While these functions are of some value in their own right, they are even more useful as examples to modify for your own purposes. The functions are general enough to be used with any table, but you have to specify table and field names (as described below) while creating a trigger.

Each of the groups of functions described below is provided as a separately-installable extension.

F.36.1. refint — Functions for Implementing Referential Integrity

\texttt{check\_primary\_key()} and \texttt{check\_foreign\_key()} are used to check foreign key constraints. (This functionality is long since superseded by the built-in foreign key mechanism, of course, but the module is still useful as an example.)

\texttt{check\_primary\_key()} checks the referencing table. To use, create a \texttt{BEFORE INSERT OR UPDATE} trigger using this function on a table referencing another table. Specify as the trigger arguments: the referencing table’s column name(s) which form the foreign key, the referenced table name, and the column names in the referenced table which form the primary/unique key. To handle multiple foreign keys, create a trigger for each reference.

\texttt{check\_foreign\_key()} checks the referenced table. To use, create a \texttt{BEFORE DELETE OR UPDATE} trigger using this function on a table referenced by other table(s). Specify as the trigger arguments: the number of referencing tables for which the function has to perform checking, the action if a referencing key is found (\texttt{cascade} — to delete the referencing row, \texttt{restrict} — to abort transaction if referencing keys exist, \texttt{setnull} — to set referencing key fields to null), the triggered table’s column names which form the primary/unique key, then the referencing table name and column names (repeated for as many referencing tables as were specified by first argument). Note that the primary/unique key columns should be marked \texttt{NOT NULL} and should have a unique index.

There are examples in \texttt{refint.example}.

F.36.2. timetravel — Functions for Implementing Time Travel

Long ago, PostgreSQL had a built-in time travel feature that kept the insert and delete times for each tuple. This can be emulated using these functions. To use these functions, you must add to a table two columns of \texttt{abstime} type to store the date when a tuple was inserted (\texttt{start\_date}) and changed/deleted (\texttt{stop\_date}):

\begin{verbatim}
CREATE TABLE mytab ( ...
    start_date abstime,
    stop_date abstime
    ...
) ;
\end{verbatim}

The columns can be named whatever you like, but in this discussion we’ll call them \texttt{start\_date} and \texttt{stop\_date}.

When a new row is inserted, \texttt{start\_date} should normally be set to current time, and \texttt{stop\_date} to \texttt{infinity}. The trigger will automatically substitute these values if the inserted data contains nulls in these columns. Generally, inserting explicit non-null data in these columns should only be done when re-loading dumped data.

Tuples with \texttt{stop\_date} equal to \texttt{infinity} are “valid now”, and can be modified. Tuples with a finite \texttt{stop\_date} cannot be modified anymore — the trigger will prevent it. (If you need to do that, you can turn off time travel as shown below.)

For a modifiable row, on update only the \texttt{stop\_date} in the tuple being updated will be changed (to current time) and a new tuple with the modified data will be inserted. \texttt{Start\_date} in this new tuple will be set to current time and \texttt{stop\_date} to \texttt{infinity}.
Appendix F. Additional Supplied Modules

A delete does not actually remove the tuple but only sets its stop_date to current time.

To query for tuples “valid now”, include stop_date = ‘infinity’ in the query’s WHERE condition. (You might wish to incorporate that in a view.) Similarly, you can query for tuples valid at any past time with suitable conditions on start_date and stop_date.

timetavel() is the general trigger function that supports this behavior. Create a BEFORE INSERT OR UPDATE OR DELETE trigger using this function on each time-traveled table. Specify two trigger arguments: the actual names of the start_date and stop_date columns. Optionally, you can specify one to three more arguments, which must refer to columns of type text. The trigger will store the name of the current user into the first of these columns during INSERT, the second column during UPDATE, and the third during DELETE.

set_timetravel() allows you to turn time-travel on or off for a table.

set_timetravel(‘mytab’, 1) will turn TT ON for table mytab. set_timetravel(‘mytab’, 0) will turn TT OFF for table mytab. In both cases the old status is reported. While TT is off, you can modify the start_date and stop_date columns freely. Note that the on/off status is local to the current database session — fresh sessions will always start out with TT ON for all tables.

get_timetravel() returns the TT state for a table without changing it.

There is an example in timetravel.example.

F.36.3. autoinc — Functions for Autoincrementing Fields

autoinc() is a trigger that stores the next value of a sequence into an integer field. This has some overlap with the built-in “serial column” feature, but it is not the same: autoinc() will override attempts to substitute a different field value during inserts, and optionally it can be used to increment the field during updates, too.

To use, create a BEFORE INSERT (or optionally BEFORE INSERT OR UPDATE) trigger using this function. Specify two trigger arguments: the name of the integer column to be modified, and the name of the sequence object that will supply values. (Actually, you can specify any number of pairs of such names, if you’d like to update more than one autoincrementing column.)

There is an example in autoinc.example.

F.36.4. insert_username — Functions for Tracking Who Changed a Table

insert_username() is a trigger that stores the current user’s name into a text field. This can be useful for tracking who last modified a particular row within a table.

To use, create a BEFORE INSERT and/or UPDATE trigger using this function. Specify a single trigger argument: the name of the text column to be modified.

There is an example in insert_username.example.

F.36.5. moddatetime — Functions for Tracking Last Modification Time

moddatetime() is a trigger that stores the current time into a timestamp field. This can be useful for tracking the last modification time of a particular row within a table.
To use, create a `BEFORE UPDATE` trigger using this function. Specify a single trigger argument: the name of the column to be modified. The column must be of type `timestamp` or `timestamp with time zone`.

There is an example in `moddatetime.example`.

### F.37. sslinfo

The `sslinfo` module provides information about the SSL certificate that the current client provided when connecting to PostgreSQL. The module is useless (most functions will return NULL) if the current connection does not use SSL.

This extension won’t build at all unless the installation was configured with `--with-openssl`.

#### F.37.1. Functions Provided

- `ssl_is_used()` returns boolean
  - Returns TRUE if current connection to server uses SSL, and FALSE otherwise.

- `ssl_version()` returns text
  - Returns the name of the protocol used for the SSL connection (e.g. SSLv2, SSLv3, or TLSv1).

- `ssl_cipher()` returns text
  - Returns the name of the cipher used for the SSL connection (e.g. DHE-RSA-AES256-SHA).

- `ssl_client_cert_present()` returns boolean
  - Returns TRUE if current client has presented a valid SSL client certificate to the server, and FALSE otherwise. (The server might or might not be configured to require a client certificate.)

- `ssl_client_serial()` returns numeric
  - Returns serial number of current client certificate. The combination of certificate serial number and certificate issuer is guaranteed to uniquely identify a certificate (but not its owner — the owner ought to regularly change their keys, and get new certificates from the issuer).

  So, if you run your own CA and allow only certificates from this CA to be accepted by the server, the serial number is the most reliable (albeit not very mnemonic) means to identify a user.

- `ssl_client_dn()` returns text
  - Returns the full subject of the current client certificate, converting character data into the current database encoding. It is assumed that if you use non-ASCII characters in the certificate names, your database is able to represent these characters, too. If your database uses the `SQL_ASCII` encoding, non-ASCII characters in the name will be represented as UTF-8 sequences.

  The result looks like `/CN=Somebody /C=Some country/O=Some organization`.

- `ssl_issuer_dn()` returns text
  - Returns the full issuer name of the current client certificate, converting character data into the current database encoding. Encoding conversions are handled the same as for `ssl_client_dn`.

  The combination of the return value of this function with the certificate serial number uniquely identifies the certificate.
This function is really useful only if you have more than one trusted CA certificate in your server’s root.crt file, or if this CA has issued some intermediate certificate authority certificates.

`ssl_client_dn_field(fieldname text) returns text`

This function returns the value of the specified field in the certificate subject, or NULL if the field is not present. Field names are string constants that are converted into ASN1 object identifiers using the OpenSSL object database. The following values are acceptable:

- commonName (alias CN)
- surname (alias SN)
- name
- givenName (alias GN)
- countryName (alias C)
- localityName (alias L)
- stateOrProvinceName (alias ST)
- organizationName (alias O)
- organizationalUnitName (alias OU)
- title
- description
- initials
- postalCode
- streetAddress
- generationQualifier
- description
- dnQualifier
- x500UniqueIdentifier
- pseudonym
- role
- emailAddress

All of these fields are optional, except commonName. It depends entirely on your CA’s policy which of them would be included and which wouldn’t. The meaning of these fields, however, is strictly defined by the X.500 and X.509 standards, so you cannot just assign arbitrary meaning to them.

`ssl_issuer_field(fieldname text) returns text`

Same as `ssl_client_dn_field`, but for the certificate issuer rather than the certificate subject.

`ssl_extension_info() returns setof record`

Provide information about extensions of client certificate: extension name, extension value, and if it is a critical extension.

F.37.2. Author

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Appendix F. Additional Supplied Modules

F.38. tablefunc

The `tablefunc` module includes various functions that return tables (that is, multiple rows). These functions are useful both in their own right and as examples of how to write C functions that return multiple rows.

F.38.1. Functions Provided

Table F-31 shows the functions provided by the `tablefunc` module.

Table F-31. tablefunc Functions

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>normal_rand(int numvals, float8 mean, float8 stddev)</td>
<td>setof float8</td>
<td>Produces a set of normally distributed random values</td>
</tr>
<tr>
<td>crosstab(text sql)</td>
<td>setof record</td>
<td>Produces a “pivot table” containing row names plus $N$ value columns, where $N$ is determined by the row type specified in the calling query</td>
</tr>
<tr>
<td>crosstabN(text sql)</td>
<td>setof table_crosstab_N</td>
<td>Produces a “pivot table” containing row names plus $N$ value columns. $crosstab2$, $crosstab3$, and $crosstab4$ are predefined, but you can create additional $crosstabN$ functions as described below</td>
</tr>
<tr>
<td>crosstab(text source_sql, text category_sql)</td>
<td>setof record</td>
<td>Produces a “pivot table” with the value columns specified by a second query</td>
</tr>
<tr>
<td>crosstab(text sql, int N)</td>
<td>setof record</td>
<td>Obsolete version of $crosstab(text)$. The parameter $N$ is now ignored, since the number of value columns is always determined by the calling query</td>
</tr>
<tr>
<td>connectby(text relname, text keyid_fld, text parent_keyid_fld [, text orderby_fld ], text start_with, int max_depth [, text branch_delim ])</td>
<td>setof record</td>
<td>Produces a representation of a hierarchical tree structure</td>
</tr>
</tbody>
</table>

F.38.1.1. normal_rand

`normal_rand(int numvals, float8 mean, float8 stddev) returns setof float8`
normal_rand produces a set of normally distributed random values (Gaussian distribution).

numvals is the number of values to be returned from the function. mean is the mean of the normal distribution of values and stddev is the standard deviation of the normal distribution of values.

For example, this call requests 1000 values with a mean of 5 and a standard deviation of 3:

test=# SELECT * FROM normal_rand(1000, 5, 3);

<table>
<thead>
<tr>
<th>normal_rand</th>
</tr>
</thead>
<tbody>
<tr>
<td>1.56556322244898</td>
</tr>
<tr>
<td>9.10040991424657</td>
</tr>
<tr>
<td>5.36957140345079</td>
</tr>
<tr>
<td>-0.369151492880995</td>
</tr>
<tr>
<td>0.283600703686639</td>
</tr>
<tr>
<td>.</td>
</tr>
<tr>
<td>.</td>
</tr>
<tr>
<td>4.82992125404908</td>
</tr>
<tr>
<td>9.71308014517282</td>
</tr>
<tr>
<td>2.49639286969028</td>
</tr>
</tbody>
</table>

(1000 rows)

F.38.1.2. crosstab(text)

crosstab(text sql)
crosstab(text sql, int N)

The crosstab function is used to produce “pivot” displays, wherein data is listed across the page rather than down. For example, we might have data like

```
row1  val11
row1  val12
row1  val13
...
row2  val21
row2  val22
row2  val23
...
```

which we wish to display like

```
row1  val11  val12  val13  ...
row2  val21  val22  val23  ...
...
```

The crosstab function takes a text parameter that is a SQL query producing raw data formatted in the first way, and produces a table formatted in the second way.

The sql parameter is a SQL statement that produces the source set of data. This statement must return one row_name column, one category column, and one value column. N is an obsolete parameter, ignored if supplied (formerly this had to match the number of output value columns, but now that is determined by the calling query).

For example, the provided query might produce a set something like:

```
row_name  cat  value
```
Appendix F. Additional Supplied Modules

---
row1  cat1  val1
row1  cat2  val2
row1  cat3  val3
row1  cat4  val4
row2  cat1  val5
row2  cat2  val6
row2  cat3  val7
row2  cat4  val8
---

The crosstab function is declared to return setof record, so the actual names and types of the output columns must be defined in the FROM clause of the calling SELECT statement, for example:

```
SELECT * FROM crosstab(‘...’) AS ct(row_name text, category_1 text, category_2 text);
```

This example produces a set something like:

```
<== value columns ==>
row_name category_1 category_2
----------+------------+------------
row1      val1      val2
row2      val5      val6
```

The FROM clause must define the output as one row_name column (of the same data type as the first result column of the SQL query) followed by N value columns (all of the same data type as the third result column of the SQL query). You can set up as many output value columns as you wish. The names of the output columns are up to you.

The crosstab function produces one output row for each consecutive group of input rows with the same row_name value. It fills the output value columns, left to right, with the value fields from these rows. If there are fewer rows in a group than there are output value columns, the extra output columns are filled with nulls; if there are more rows, the extra input rows are skipped.

In practice the SQL query should always specify ORDER BY 1,2 to ensure that the input rows are properly ordered, that is, values with the same row_name are brought together and correctly ordered within the row. Notice that crosstab itself does not pay any attention to the second column of the query result; it’s just there to be ordered by, to control the order in which the third-column values appear across the page.

Here is a complete example:

```
CREATE TABLE ct(id SERIAL, rowid TEXT, attribute TEXT, value TEXT);
INSERT INTO ct(rowid, attribute, value) VALUES('test1','att1','val1');
INSERT INTO ct(rowid, attribute, value) VALUES('test1','att2','val2');
INSERT INTO ct(rowid, attribute, value) VALUES('test1','att3','val3');
INSERT INTO ct(rowid, attribute, value) VALUES('test1','att4','val4');
INSERT INTO ct(rowid, attribute, value) VALUES('test2','att1','val5');
INSERT INTO ct(rowid, attribute, value) VALUES('test2','att2','val6');
INSERT INTO ct(rowid, attribute, value) VALUES('test2','att3','val7');
INSERT INTO ct(rowid, attribute, value) VALUES('test2','att4','val8');

SELECT *
FROM crosstab(
  'select rowid, attribute, value
  from ct
```
where attribute = "att2" or attribute = "att3"
order by 1,2'
AS ct(row_name text, category_1 text, category_2 text, category_3 text);

row_name | category_1 | category_2 | category_3
----------+------------+------------+------------
test1    | val2       | val3       |
test2    | val6       | val7       |
----------+------------+------------+------------
(2 rows)

You can avoid always having to write out a FROM clause to define the output columns, by setting up a custom crosstab function that has the desired output row type wired into its definition. This is described in the next section. Another possibility is to embed the required FROM clause in a view definition.

Note: See also the \crosstabview command in psql, which provides functionality similar to crosstab().

F.38.1.3. crosstabcN (text)

crosstabcN(text sql)

The crosstabcN functions are examples of how to set up custom wrappers for the general crosstab function, so that you need not write out column names and types in the calling SELECT query. The tablefunc module includes crosstab2, crosstab3, and crosstab4, whose output row types are defined as

CREATE TYPE tablefunc_crosstab_N AS {
  row_name TEXT,
  category_1 TEXT,
  category_2 TEXT,
  .
  .
  .
  category_N TEXT
};

Thus, these functions can be used directly when the input query produces row_name and value columns of type text, and you want 2, 3, or 4 output values columns. In all other ways they behave exactly as described above for the general crosstab function.

For instance, the example given in the previous section would also work as

```
SELECT *
FROM crosstab3(
  'select rowid, attribute, value
   from ct
   where attribute = "att2" or attribute = "att3"
   order by 1,2');
```
These functions are provided mostly for illustration purposes. You can create your own return types and functions based on the underlying `crosstab()` function. There are two ways to do it:

- Create a composite type describing the desired output columns, similar to the examples in `contrib/tablefunc/tablefunc--1.0.sql`. Then define a unique function name accepting one text parameter and returning `setof my_type_name`, but linking to the same underlying `crosstab` C function. For example, if your source data produces row names that are `text`, and values that are `float8`, and you want 5 value columns:

```sql
CREATE TYPE my_crosstab_float8_5_cols AS (
    my_row_name text,
    my_category_1 float8,
    my_category_2 float8,
    my_category_3 float8,
    my_category_4 float8,
    my_category_5 float8
);

CREATE OR REPLACE FUNCTION crosstab_float8_5_cols(text)
RETURNS setof my_crosstab_float8_5_cols
AS '$libdir/tablefunc','crosstab' LANGUAGE C STABLE STRICT;
```

- Use `OUT` parameters to define the return type implicitly. The same example could also be done this way:

```sql
CREATE OR REPLACE FUNCTION crosstab_float8_5_cols(
    IN text,
    OUT my_row_name text,
    OUT my_category_1 float8,
    OUT my_category_2 float8,
    OUT my_category_3 float8,
    OUT my_category_4 float8,
    OUT my_category_5 float8)
RETURNS setof record
AS '$libdir/tablefunc','crosstab' LANGUAGE C STABLE STRICT;
```

\[\text{F.38.1.4. crosstab(text, text)}\]

crosstab(text source_sql, text category_sql)

The main limitation of the single-parameter form of `crosstab` is that it treats all values in a group alike, inserting each value into the first available column. If you want the value columns to correspond to specific categories of data, and some groups might not have data for some of the categories, that doesn’t work well. The two-parameter form of `crosstab` handles this case by providing an explicit list of the categories corresponding to the output columns.

`source_sql` is a SQL statement that produces the source set of data. This statement must return one `row_name` column, one `category` column, and one `value` column. It may also have one or more “extra” columns. The `row_name` column must be first. The `category` and `value` columns must be the last two columns, in that order. Any columns between `row_name` and `category` are treated as “extra”. The “extra” columns are expected to be the same for all rows with the same `row_name` value.

For example, `source_sql` might produce a set something like:

```sql
SELECT row_name, extra_col, cat, value FROM foo ORDER BY 1;
```
category_sql is a SQL statement that produces the set of categories. This statement must return only one column. It must produce at least one row, or an error will be generated. Also, it must not produce duplicate values, or an error will be generated. category_sql might be something like:

```
SELECT DISTINCT cat FROM foo ORDER BY 1;
```

```
cat
---
cat1
cat2
cat3
cat4
```

The crosstab function is declared to return setof record, so the actual names and types of the output columns must be defined in the FROM clause of the calling SELECT statement, for example:

```
SELECT * FROM crosstab('...', '...')
AS ct(row_name text, extra text, cat1 text, cat2 text, cat3 text, cat4 text);
```

This will produce a result something like:

```
television channel -------------------------------------
row_name extra cat1 cat2 cat3 cat4
---------+-------+------+------+------+------
row1 extra1 val1 val2 val4
row2 extra2 val5 val6 val7 val8
```

The FROM clause must define the proper number of output columns of the proper data types. If there are N columns in the source_sql query’s result, the first N-2 of them must match up with the first N-2 output columns. The remaining output columns must have the type of the last column of the source_sql query’s result, and there must be exactly as many of them as there are rows in the category_sql query’s result.

The crosstab function produces one output row for each consecutive group of input rows with the same row_name value. The output row_name column, plus any “extra” columns, are copied from the first row of the group. The output value columns are filled with the value fields from rows having matching category values. If a row’s category does not match any output of the category_sql query, its value is ignored. Output columns whose matching category is not present in any input row of the group are filled with nulls.
In practice the `source_sql` query should always specify `ORDER BY 1` to ensure that values with the same `row_name` are brought together. However, ordering of the categories within a group is not important. Also, it is essential to be sure that the order of the `category_sql` query’s output matches the specified output column order.

Here are two complete examples:

```sql
create table sales(year int, month int, qty int);
insert into sales values(2007, 1, 1000);
insert into sales values(2007, 2, 1500);
insert into sales values(2007, 7, 500);
insert into sales values(2007, 11, 1500);
insert into sales values(2007, 12, 2000);
insert into sales values(2008, 1, 1000);

select * from crosstab(
    'select year, month, qty from sales order by 1',
    'select m from generate_series(1,12) m'
) as (year int,
    "Jan" int,
    "Feb" int,
    "Mar" int,
    "Apr" int,
    "May" int,
    "Jun" int,
    "Jul" int,
    "Aug" int,
    "Sep" int,
    "Oct" int,
    "Nov" int,
    "Dec" int
);

<table>
<thead>
<tr>
<th>year</th>
<th>Jan</th>
<th>Feb</th>
<th>Mar</th>
<th>Apr</th>
<th>May</th>
<th>Jun</th>
<th>Jul</th>
<th>Aug</th>
<th>Sep</th>
<th>Oct</th>
<th>Nov</th>
<th>Dec</th>
</tr>
</thead>
<tbody>
<tr>
<td>2007</td>
<td>1000</td>
<td>1500</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>500</td>
<td></td>
<td></td>
<td>1500</td>
<td>2000</td>
<td></td>
</tr>
<tr>
<td>2008</td>
<td>1000</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
```

CREATE TABLE cth(rowid text, rowdt timestamp, attribute text, val text);
insert into cth values('test1','01 March 2003','temperature','42');
insert into cth values('test1','01 March 2003','test_result','PASS');
insert into cth values('test1','01 March 2003','volts','2.6987');
insert into cth values('test2','02 March 2003','temperature','53');
insert into cth values('test2','02 March 2003','test_result','FAIL');
insert into cth values('test2','02 March 2003','test_startdate','01 March 2003');
insert into cth values('test2','02 March 2003','volts','3.1234');

SELECT * FROM crosstab
(    'SELECT rowid, rowdt, attribute, val FROM cth ORDER BY 1',
    'SELECT DISTINCT attribute FROM cth ORDER BY 1'
) AS
(    rowid text,
    rowdt timestamp,
    temperature int4,
);  ```
Appendix F: Additional Supplied Modules

test_result text,
test_startdate timestamp,
vols float8
);

rowid | rowdt | temperature | test_result | test_startdate | volts
-----+--------------------------+-------------+-------------+--------------------------+--------
test1 | Sat Mar 01 00:00:00 2003 | 42 | PASS | | 2.6987

test2 | Sun Mar 02 00:00:00 2003 | 53 | FAIL | Sat Mar 01 00:00:00 2003 | 3.1234

(2 rows)

You can create predefined functions to avoid having to write out the result column names and types in each query. See the examples in the previous section. The underlying C function for this form of crosstab is named crosstab_hash.

F.38.1.5. connectby

connectby(text relname, text keyid_fld, text parent_keyid_fld
[, text orderby_fld ], text start_with, int max_depth
[, text branch_delim ])

The connectby function produces a display of hierarchical data that is stored in a table. The table must have a key field that uniquely identifies rows, and a parent-key field that references the parent (if any) of each row. connectby can display the sub-tree descending from any row.

Table F-32 explains the parameters.

Table F-32. connectby Parameters

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>relname</td>
<td>Name of the source relation</td>
</tr>
<tr>
<td>keyid_fld</td>
<td>Name of the key field</td>
</tr>
<tr>
<td>parent_keyid_fld</td>
<td>Name of the parent-key field</td>
</tr>
<tr>
<td>orderby_fld</td>
<td>Name of the field to order siblings by (optional)</td>
</tr>
<tr>
<td>start_with</td>
<td>Key value of the row to start at</td>
</tr>
<tr>
<td>max_depth</td>
<td>Maximum depth to descend to, or zero for unlimited depth</td>
</tr>
<tr>
<td>branch_delim</td>
<td>String to separate keys with in branch output (optional)</td>
</tr>
</tbody>
</table>

The key and parent-key fields can be any data type, but they must be the same type. Note that the start_with value must be entered as a text string, regardless of the type of the key field.

The connectby function is declared to return setof record, so the actual names and types of the output columns must be defined in the FROM clause of the calling SELECT statement, for example:

```
AS t(keyid text, parent_keyid text, level int, branch text, pos int);
```

The first two output columns are used for the current row’s key and its parent row’s key; they must match the type of the table’s key field. The third output column is the depth in the tree and must be of type integer. If a branch_delim parameter was given, the next output column is the branch display and must be of type text. Finally, if an orderby_fld parameter was given, the last output
column is a serial number, and must be of type integer.

The “branch” output column shows the path of keys taken to reach the current row. The keys are separated by the specified branch_delim string. If no branch display is wanted, omit both the branch_delim parameter and the branch column in the output column list.

If the ordering of siblings of the same parent is important, include the orderby_fld parameter to specify which field to order siblings by. This field can be of any sortable data type. The output column list must include a final integer serial-number column, if and only if orderby_fld is specified.

The parameters representing table and field names are copied as-is into the SQL queries that connectby generates internally. Therefore, include double quotes if the names are mixed-case or contain special characters. You may also need to schema-qualify the table name.

In large tables, performance will be poor unless there is an index on the parent-key field.

It is important that the branch_delim string not appear in any key values, else connectby may incorrectly report an infinite-recursion error. Note that if branch_delim is not provided, a default value of ~ is used for recursion detection purposes.

Here is an example:

```sql
CREATE TABLE connectby_tree(keyid text, parent_keyid text, pos int);

INSERT INTO connectby_tree VALUES('row1',NULL, 0);
INSERT INTO connectby_tree VALUES('row2','row1', 0);
INSERT INTO connectby_tree VALUES('row3','row1', 0);
INSERT INTO connectby_tree VALUES('row4','row2', 1);
INSERT INTO connectby_tree VALUES('row5','row2', 0);
INSERT INTO connectby_tree VALUES('row6','row4', 0);
INSERT INTO connectby_tree VALUES('row7','row3', 0);
INSERT INTO connectby_tree VALUES('row8','row6', 0);
INSERT INTO connectby_tree VALUES('row9','row5', 0);

-- with branch, without orderby_fld (order of results is not guaranteed)
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'row2', 0, '~')
AS t(keyid text, parent_keyid text, level int, branch text);

keyid | parent_keyid | level | branch
-------+--------------+-------+---------------------
row2 | | 0 | row2
row4 | row2 | 1 | row2~row4
row6 | row4 | 2 | row2~row4~row6
row8 | row6 | 3 | row2~row4~row6~row8
row5 | row2 | 1 | row2~row5
row9 | row5 | 2 | row2~row5~row9

(6 rows)

-- without branch, without orderby_fld (order of results is not guaranteed)
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'row2', 0)
AS t(keyid text, parent_keyid text, level int);

keyid | parent_keyid | level
-------+--------------+-------
row2 | | 0
row4 | row2 | 1
row6 | row4 | 2
row8 | row6 | 3
row5 | row2 | 1
row9 | row5 | 2

(6 rows)
```
-- with branch, with orderby_fld (notice that row5 comes before row4)

```sql
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'pos', 'row2', 0, '~-')
AS t(keyid text, parent_keyid text, level int, branch text, pos int);
```

<table>
<thead>
<tr>
<th>keyid</th>
<th>parent_keyid</th>
<th>level</th>
<th>branch</th>
<th>pos</th>
</tr>
</thead>
<tbody>
<tr>
<td>row2</td>
<td></td>
<td>0</td>
<td>row2</td>
<td>1</td>
</tr>
<tr>
<td>row5</td>
<td>row2</td>
<td>1</td>
<td>row2~row5</td>
<td>2</td>
</tr>
<tr>
<td>row9</td>
<td>row5</td>
<td>2</td>
<td>row2<del>row5</del>row9</td>
<td>3</td>
</tr>
<tr>
<td>row4</td>
<td>row2</td>
<td>1</td>
<td>row2~row4</td>
<td>4</td>
</tr>
<tr>
<td>row6</td>
<td>row4</td>
<td>2</td>
<td>row2<del>row4</del>row6</td>
<td>5</td>
</tr>
<tr>
<td>row8</td>
<td>row6</td>
<td>3</td>
<td>row2<del>row4</del>row6~row8</td>
<td>6</td>
</tr>
</tbody>
</table>

(6 rows)

-- without branch, with orderby_fld (notice that row5 comes before row4)

```sql
SELECT * FROM connectby('connectby_tree', 'keyid', 'parent_keyid', 'pos', 'row2', 0)
AS t(keyid text, parent_keyid text, level int, pos int);
```

<table>
<thead>
<tr>
<th>keyid</th>
<th>parent_keyid</th>
<th>level</th>
<th>pos</th>
</tr>
</thead>
<tbody>
<tr>
<td>row2</td>
<td></td>
<td>0</td>
<td>1</td>
</tr>
<tr>
<td>row5</td>
<td>row2</td>
<td>1</td>
<td>2</td>
</tr>
<tr>
<td>row9</td>
<td>row5</td>
<td>2</td>
<td>3</td>
</tr>
<tr>
<td>row4</td>
<td>row2</td>
<td>1</td>
<td>4</td>
</tr>
<tr>
<td>row6</td>
<td>row4</td>
<td>2</td>
<td>5</td>
</tr>
<tr>
<td>row8</td>
<td>row6</td>
<td>3</td>
<td>6</td>
</tr>
</tbody>
</table>

(6 rows)

F.38.2. Author

Joe Conway

F.39. tcn

The tcn module provides a trigger function that notifies listeners of changes to any table on which it is attached. It must be used as an `AFTER trigger FOR EACH ROW`.

Only one parameter may be supplied to the function in a `CREATE TRIGGER` statement, and that is optional. If supplied it will be used for the channel name for the notifications. If omitted `tcn` will be used for the channel name.

The payload of the notifications consists of the table name, a letter to indicate which type of operation was performed, and column name/value pairs for primary key columns. Each part is separated from the next by a comma. For ease of parsing using regular expressions, table and column names are always wrapped in double quotes, and data values are always wrapped in single quotes. Embedded quotes are doubled.

A brief example of using the extension follows.

```sql
test=# create table tcndata
```

```sql
test=# ( test(# a int not null,
```
test(#   b date not null,
test(#   c text,
test(#   primary key (a, b)
test(# );
CREATE TABLE
test=# create trigger tcndata_tcn_trigger
  test-#  after insert or update or delete on tcndata
  test-#   for each row execute procedure triggered_change_notification();
CREATE TRIGGER
test=# listen tcn;
LISTEN
test=# insert into tcndata values (1, date '2012-12-22', 'one'),
test-#   (1, date '2012-12-23', 'another'),
test-#   (2, date '2012-12-23', 'two');
INSERT 0 3
Asynchronous notification "tcn" with payload ""tcndata",I,"a"='1',"b"='2012-12-22'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",I,"a"='1',"b"='2012-12-23'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",I,"a"='2',"b"='2012-12-23'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",U,"a"='1',"b"='2012-12-22'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",U,"a"='1',"b"='2012-12-23'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",D,"a"='1',"b"='2012-12-22'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",D,"a"='1',"b"='2012-12-23'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",D,"a"='2',"b"='2012-12-23'" received from server process with PID 22770.
update tcndata set c = 'uno' where a = 1;
UPDATE 2
Asynchronous notification "tcn" with payload ""tcndata",U,"a"='1',"b"='2012-12-22'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",U,"a"='1',"b"='2012-12-23'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",D,"a"='1',"b"='2012-12-22'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",D,"a"='1',"b"='2012-12-23'" received from server process with PID 22770.
Asynchronous notification "tcn" with payload ""tcndata",D,"a"='2',"b"='2012-12-23'" received from server process with PID 22770.

F.40. test_decoding

test_decoding is an example of a logical decoding output plugin. It doesn’t do anything especially useful, but can serve as a starting point for developing your own output plugin.
test_decoding receives WAL through the logical decoding mechanism and decodes it into text representations of the operations performed.

Typical output from this plugin, used over the SQL logical decoding interface, might be:

postgres=# SELECT * FROM pg_logical_slot_get_changes('test_slot', NULL, NULL, 'include-xids', '0');
location | xid | data
----------+-----+--------------------------------------------------
0/16D30F8 | 691 | BEGIN
0/16D32A0 | 691 | table public.data: INSERT: id[int4]:2 data[text]:'arg'
0/16D32A0 | 691 | table public.data: INSERT: id[int4]:3 data[text]:'demo'
0/16D32A0 | 691 | COMMIT
0/16D32D8 | 692 | BEGIN
0/16D3398 | 692 | table public.data: DELETE: id[int4]:2
0/16D3398 | 692 | table public.data: DELETE: id[int4]:3
0/16D3398 | 692 | COMMIT
(8 rows)
Appendix F. Additional Supplied Modules

F.41. tsearch2

The tsearch2 module provides backwards-compatible text search functionality for applications that used tsearch2 before text searching was integrated into core PostgreSQL in release 8.3.

F.41.1. Portability Issues

Although the built-in text search features were based on tsearch2 and are largely similar to it, there are numerous small differences that will create portability issues for existing applications:

- Some functions’ names were changed, for example rank to ts_rank. The replacement tsearch2 module provides aliases having the old names.
- The built-in text search data types and functions all exist within the system schema pg_catalog. In an installation using tsearch2, these objects would usually have been in the public schema, though some users chose to place them in a separate schema of their own. Explicitly schema-qualified references to the objects will therefore fail in either case. The replacement tsearch2 module provides alias objects that are stored in public (or another schema if necessary) so that such references will still work.
- There is no concept of a “current parser” or “current dictionary” in the built-in text search features, only of a current search configuration (set by the default_text_search_config parameter). While the current parser and current dictionary were used only by functions intended for debugging, this might still pose a porting obstacle in some cases. The replacement tsearch2 module emulates these additional state variables and provides backwards-compatible functions for setting and retrieving them.

There are some issues that are not addressed by the replacement tsearch2 module, and will therefore require application code changes in any case:

- The old tsearch2 trigger function allowed items in its argument list to be names of functions to be invoked on the text data before it was converted to tsvector format. This was removed as being a security hole, since it was not possible to guarantee that the function invoked was the one intended. The recommended approach if the data must be massaged before being indexed is to write a custom trigger that does the work for itself.
- Text search configuration information has been moved into core system catalogs that are noticeably different from the tables used by tsearch2. Any applications that examined or modified those tables will need adjustment.
- If an application used any custom text search configurations, those will need to be set up in the core catalogs using the new text search configuration SQL commands. The replacement tsearch2 module offers a little bit of support for this by making it possible to load an old set of tsearch2 configuration tables into PostgreSQL 8.3. (Without the module, it is not possible to load the configuration data because values in the regprocedure columns cannot be resolved to functions.) While those configuration tables won’t actually do anything, at least their contents will be available to be consulted while setting up an equivalent custom configuration in 8.3.
- The old reset_tsearch() and get_covers() functions are not supported.
- The replacement tsearch2 module does not define any alias operators, relying entirely on the built-in ones. This would only pose an issue if an application used explicitly schema-qualified operator names, which is very uncommon.
F.41.2. Converting a pre-8.3 Installation

The recommended way to update a pre-8.3 installation that uses tsearch2 is:

1. Make a dump from the old installation in the usual way, but be sure not to use \(-c (--clean)\) option of pg_dump or pg_dumpall.
2. In the new installation, create empty database(s) and install the replacement tsearch2 module into each database that will use text search. This must be done before loading the dump data! If your old installation had the tsearch2 objects in a schema other than public, be sure to adjust the CREATE EXTENSION command so that the replacement objects are created in that same schema.
3. Load the dump data. There will be quite a few errors reported due to failure to recreate the original tsearch2 objects. These errors can be ignored, but this means you cannot restore the dump in a single transaction (eg, you cannot use pg_restore’s \(-1\) switch).
4. Examine the contents of the restored tsearch2 configuration tables (pg_ts_cfg and so on), and create equivalent built-in text search configurations as needed. You may drop the old configuration tables once you’ve extracted all the useful information from them.
5. Test your application.

At a later time you may wish to rename application references to the alias text search objects, so that you can eventually uninstall the replacement tsearch2 module.

F.41.3. References

Tsearch2 Development Site http://www.sai.msu.su/~megera/postgres/gist/tsearch/V2/

F.42. tsm_system_rows

The tsm_system_rows module provides the table sampling method SYSTEM_ROWS, which can be used in the TABLESAMPLE clause of a SELECT command.

This table sampling method accepts a single integer argument that is the maximum number of rows to read. The resulting sample will always contain exactly that many rows, unless the table does not contain enough rows, in which case the whole table is selected.

Like the built-in SYSTEM sampling method, SYSTEM_ROWS performs block-level sampling, so that the sample is not completely random but may be subject to clustering effects, especially if only a small number of rows are requested.

SYSTEM_ROWS does not support the REPEATABLE clause.

F.42.1. Examples

Here is an example of selecting a sample of a table with SYSTEM_ROWS. First install the extension:

```
CREATE EXTENSION tsm_system_rows;
```

Then you can use it in a SELECT command, for instance:

```
SELECT * FROM my_table TABLESAMPLE SYSTEM_ROWS(100);
```
This command will return a sample of 100 rows from the table my_table (unless the table does not have 100 visible rows, in which case all its rows are returned).

F.43. tsm_system_time

The tsm_system_time module provides the table sampling method SYSTEM_TIME, which can be used in the TABLESAMPLE clause of a SELECT command.

This table sampling method accepts a single floating-point argument that is the maximum number of milliseconds to spend reading the table. This gives you direct control over how long the query takes, at the price that the size of the sample becomes hard to predict. The resulting sample will contain as many rows as could be read in the specified time, unless the whole table has been read first.

Like the built-in SYSTEM sampling method, SYSTEM_TIME performs block-level sampling, so that the sample is not completely random but may be subject to clustering effects, especially if only a small number of rows are selected.

SYSTEM_TIME does not support the REPEATABLE clause.

F.43.1. Examples

Here is an example of selecting a sample of a table with SYSTEM_TIME. First install the extension:

CREATE EXTENSION tsm_system_time;

Then you can use it in a SELECT command, for instance:

SELECT * FROM my_table TABLESAMPLE SYSTEM_TIME(1000);

This command will return as large a sample of my_table as it can read in 1 second (1000 milliseconds). Of course, if the whole table can be read in under 1 second, all its rows will be returned.

F.44. unaccent

unaccent is a text search dictionary that removes accents (diacritic signs) from lexemes. It’s a filtering dictionary, which means its output is always passed to the next dictionary (if any), unlike the normal behavior of dictionaries. This allows accent-insensitive processing for full text search.

The current implementation of unaccent cannot be used as a normalizing dictionary for the thesaurus dictionary.

F.44.1. Configuration

An unaccent dictionary accepts the following options:
RULES is the base name of the file containing the list of translation rules. This file must be stored in $SHAREDIR/tsearch_data/ (where $SHAREDIR means the PostgreSQL installation’s shared-data directory). Its name must end in .rules (which is not to be included in the RULES parameter).

The rules file has the following format:

- Each line represents one translation rule, consisting of a character with accent followed by a character without accent. The first is translated into the second. For example,

  A \ A
  A \ A
  A \ A
  Ä \ A
  Å \ A
  Â \ A
  E \ AE

  The two characters must be separated by whitespace, and any leading or trailing whitespace on a line is ignored.

- Alternatively, if only one character is given on a line, instances of that character are deleted; this is useful in languages where accents are represented by separate characters.

- Actually, each “character” can be any string not containing whitespace, so unaccent dictionaries could be used for other sorts of substring substitutions besides diacritic removal.

- As with other PostgreSQL text search configuration files, the rules file must be stored in UTF-8 encoding. The data is automatically translated into the current database’s encoding when loaded. Any lines containing untranslatable characters are silently ignored, so that rules files can contain rules that are not applicable in the current encoding.

A more complete example, which is directly useful for most European languages, can be found in unaccent.rules, which is installed in $SHAREDIR/tsearch_data/ when the unaccent module is installed. This rules file translates characters with accents to the same characters without accents, and it also expands ligatures into the equivalent series of simple characters (for example, Ä to AE).

F.44.2. Usage

Installing the unaccent extension creates a text search template unaccent and a dictionary unaccent based on it. The unaccent dictionary has the default parameter setting RULES='unaccent', which makes it immediately usable with the standard unaccent.rules file. If you wish, you can alter the parameter, for example

```
mydb=# ALTER TEXT SEARCH DICTIONARY unaccent (RULES='my_rules');
```

or create new dictionaries based on the template.

To test the dictionary, you can try:

```
mydb=# select ts_lexize('unaccent','Hôtel');
ts_lexize
-----------
(1 row)
```

Here is an example showing how to insert the unaccent dictionary into a text search configuration:
F.44.3. Functions

The `unaccent()` function removes accents (diacritic signs) from a given string. Basically, it’s a wrapper around `unaccent`-type dictionaries, but it can be used outside normal text search contexts.

```
unaccent([dictionary regdictionary, ] string text) returns text
```

If the `dictionary` argument is omitted, the text search dictionary named `unaccent` and appearing in the same schema as the `unaccent()` function itself is used.

For example:

```
SELECT unaccent('unaccent', 'Hôtel');
SELECT unaccent('Hôtel');
```

F.45. uuid-osspp

The `uuid-ossp` module provides functions to generate universally unique identifiers (UUIDs) using one of several standard algorithms. There are also functions to produce certain special UUID constants.

F.45.1. uuid-osspp Functions

Table F-33 shows the functions available to generate UUIDs. The relevant standards ITU-T Rec. X.667, ISO/IEC 9834-8:2005, and RFC 4122 specify four algorithms for generating UUIDs, identi-
Appendix F. Additional Supplied Modules

Fied by the version numbers 1, 3, 4, and 5. (There is no version 2 algorithm.) Each of these algorithms could be suitable for a different set of applications.

Table F-33. Functions for UUID Generation

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>uuid_generate_v1()</td>
<td>This function generates a version 1 UUID. This involves the MAC address of the computer and a time stamp. Note that UUIDs of this kind reveal the identity of the computer that created the identifier and the time at which it did so, which might make it unsuitable for certain security-sensitive applications.</td>
</tr>
<tr>
<td>uuid_generate_v1mc()</td>
<td>This function generates a version 1 UUID but uses a random multicast MAC address instead of the real MAC address of the computer.</td>
</tr>
<tr>
<td>uuid_generate_v3(namespace uuid, name text)</td>
<td>This function generates a version 3 UUID in the given namespace using the specified input name. The namespace should be one of the special constants produced by the uuid_ns_*() functions shown in Table F-34. (It could be any UUID in theory.) The name is an identifier in the selected namespace. For example: <code>SELECT uuid_generate_v3(uuid_ns_url(), 'http://www.postgresql.org');</code> The name parameter will be MD5-hashed, so the cleartext cannot be derived from the generated UUID. The generation of UUIDs by this method has no random or environment-dependent element and is therefore reproducible.</td>
</tr>
<tr>
<td>uuid_generate_v4()</td>
<td>This function generates a version 4 UUID, which is derived entirely from random numbers.</td>
</tr>
<tr>
<td>uuid_generate_v5(namespace uuid, name text)</td>
<td>This function generates a version 5 UUID, which works like a version 3 UUID except that SHA-1 is used as a hashing method. Version 5 should be preferred over version 3 because SHA-1 is thought to be more secure than MD5.</td>
</tr>
</tbody>
</table>

Table F-34. Functions Returning UUID Constants

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>uuid_nil()</td>
<td>A “nil” UUID constant, which does not occur as a real UUID.</td>
</tr>
<tr>
<td>uuid_ns_dns()</td>
<td>Constant designating the DNS namespace for UUIDs.</td>
</tr>
<tr>
<td>uuid_ns_url()</td>
<td>Constant designating the URL namespace for UUIDs.</td>
</tr>
</tbody>
</table>
Appendix F. Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>uuid_ns_oid()</td>
<td>Constant designating the ISO object identifier (OID) namespace for UUIDs.</td>
</tr>
<tr>
<td>uuid_ns_x500()</td>
<td>Constant designating the X.500 distinguished name (DN) namespace for UUIDs.</td>
</tr>
</tbody>
</table>

F.45.2. Building uuid-ossp

Historically this module depended on the OSSP UUID library, which accounts for the module’s name. While the OSSP UUID library can still be found at http://www.ossp.org/pkg/lib/uuid/, it is not well maintained, and is becoming increasingly difficult to port to newer platforms. `uuid-ossp` can now be built without the OSSP library on some platforms. On FreeBSD, NetBSD, and some other BSD-derived platforms, suitable UUID creation functions are included in the core `libc` library. On Linux, OS X, and some other platforms, suitable functions are provided in the `libuuid` library, which originally came from the e2fsprogs project (though on modern Linux it is considered part of `util-linux-ng`). When invoking `configure`, specify `--with-uuid=bsd` to use the BSD functions, or `--with-uuid=e2fs` to use `e2fsprogs’ libuuid`, or `--with-uuid=ossp` to use the OSSP UUID library. More than one of these libraries might be available on a particular machine, so `configure` does not automatically choose one.

**Note:** If you only need randomly-generated (version 4) UUIDs, consider using the `gen_random_uuid()` function from the `pgcrypto` module instead.

F.45.3. Author

Peter Eisentraut <peter_e@gmx.net>

F.46. xml2

The `xml2` module provides XPath querying and XSLT functionality.

F.46.1. Deprecation Notice

From PostgreSQL 8.3 on, there is XML-related functionality based on the SQL/XML standard in the core server. That functionality covers XML syntax checking and XPath queries, which is what this module does, and more, but the API is not at all compatible. It is planned that this module will be removed in a future version of PostgreSQL in favor of the newer standard API, so you are encouraged to try converting your applications. If you find that some of the functionality of this module is not available in an adequate form with the newer API, please explain your issue to <pgsql-hackers@lists.postgresql.org> so that the deficiency can be addressed.
## F.46.2. Description of Functions

Table F-35 shows the functions provided by this module. These functions provide straightforward XML parsing and XPath queries. All arguments are of type `text`, so for brevity that is not shown.

**Table F-35. Functions**

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>xml_is_well_formed(document)</code></td>
<td><code>bool</code></td>
<td>This parses the document text in its parameter and returns true if the document is well-formed XML. <em>(Note: before PostgreSQL 8.2, this function was called <code>xml_valid()</code>. That is the wrong name since validity and well-formedness have different meanings in XML. The old name is still available, but is deprecated.)</em></td>
</tr>
<tr>
<td><code>xpath_string(document, query)</code></td>
<td><code>text</code></td>
<td>These functions evaluate the XPath query on the supplied document, and cast the result to the specified type.</td>
</tr>
<tr>
<td><code>xpath_number(document, query)</code></td>
<td><code>float4</code></td>
<td></td>
</tr>
<tr>
<td><code>xpath_bool(document, query)</code></td>
<td><code>bool</code></td>
<td></td>
</tr>
</tbody>
</table>
| `xpath_nodeset(document, query, toptag, itemtag)` | `text`  | This evaluates query on document and wraps the result in XML tags. If the result is multivalued, the output will look like:  
`<toptag>`  
`<itemtag>Value 1 which could be an XML fragment  
`<itemtag>`Value 2....`</itemtag>`  
`</toptag>`  
If either `toptag` or `itemtag` is an empty string, the relevant tag is omitted. |
| `xpath_nodeset(document, query)`              | `text`  | Like `xpath_nodeset(document, query, toptag, itemtag)` but result omits both tags. |
| `xpath_nodeset(document, query, itemtag)`     | `text`  | Like `xpath_nodeset(document, query, toptag, itemtag)` but result omits toptag. |
| `xpath_list(document, query, separator)`     | `text`  | This function returns multiple values separated by the specified separator, for example `Value 1,Value 2,Value 3 if separator is ,`. |
### Appendix F: Additional Supplied Modules

<table>
<thead>
<tr>
<th>Function</th>
<th>Returns</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>xpath_list(document, query)</td>
<td>text</td>
<td>This is a wrapper for the above function that uses , as the separator.</td>
</tr>
</tbody>
</table>

#### F.46.3. xpath_table

`xpath_table(text key, text document, text relation, text xpaths, text criteria) returns setof record`

`xpath_table` is a table function that evaluates a set of XPath queries on each of a set of documents and returns the results as a table. The primary key field from the original document table is returned as the first column of the result so that the result set can readily be used in joins. The parameters are described in Table F-36.

**Table F-36. xpath_table Parameters**

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>key</td>
<td>the name of the “key” field — this is just a field to be used as the first column of the output table, i.e., it identifies the record from which each output row came (see note below about multiple values)</td>
</tr>
<tr>
<td>document</td>
<td>the name of the field containing the XML document</td>
</tr>
<tr>
<td>relation</td>
<td>the name of the table or view containing the documents</td>
</tr>
<tr>
<td>xpaths</td>
<td>one or more XPath expressions, separated by</td>
</tr>
<tr>
<td>criteria</td>
<td>the contents of the WHERE clause. This cannot be omitted, so use true or 1=1 if you want to process all the rows in the relation</td>
</tr>
</tbody>
</table>

These parameters (except the XPath strings) are just substituted into a plain SQL SELECT statement, so you have some flexibility — the statement is

```
SELECT <key>, <document> FROM <relation> WHERE <criteria>
```

so those parameters can be *anything* valid in those particular locations. The result from this SELECT needs to return exactly two columns (which it will unless you try to list multiple fields for key or document). Beware that this simplistic approach requires that you validate any user-supplied values to avoid SQL injection attacks.

The function has to be used in a `FROM` expression, with an `AS` clause to specify the output columns; for example

```
SELECT * FROM
xpath_table('article_id',
    'article_xml',
    'articles',
    '/article/author|/article/pages|/article/title',
    'date_entered > "2003-01-01" ')
AS t(article_id integer, author text, page_count integer, title text);
```
The AS clause defines the names and types of the columns in the output table. The first is the “key” field and the rest correspond to the XPath queries. If there are more XPath queries than result columns, the extra queries will be ignored. If there are more result columns than XPath queries, the extra columns will be NULL.

Notice that this example defines the page_count result column as an integer. The function deals internally with string representations, so when you say you want an integer in the output, it will take the string representation of the XPath result and use PostgreSQL input functions to transform it into an integer (or whatever type the AS clause requests). An error will result if it can’t do this — for example if the result is empty — so you may wish to just stick to text as the column type if you think your data has any problems.

The calling SELECT statement doesn’t necessarily have to be just SELECT * — it can reference the output columns by name or join them to other tables. The function produces a virtual table with which you can perform any operation you wish (e.g. aggregation, joining, sorting etc). So we could also have:

```sql
SELECT t.title, p.fullname, p.email
FROM xpath_table('article_id', 'article_xml', 'articles',
    '/article/title|/article/author/@id',
    'xpath_string(article_xml,"/article/@date") > "2003-03-20" ') AS t(article_id integer, title text, author_id integer),
    tblPeopleInfo AS p
WHERE t.author_id = p.person_id;
```

as a more complicated example. Of course, you could wrap all of this in a view for convenience.

F.46.3.1. Multivalued Results

The xpath_table function assumes that the results of each XPath query might be multivalued, so the number of rows returned by the function may not be the same as the number of input documents. The first row returned contains the first result from each query, the second row the second result from each query. If one of the queries has fewer values than the others, null values will be returned instead.

In some cases, a user will know that a given XPath query will return only a single result (perhaps a unique document identifier) — if used alongside an XPath query returning multiple results, the single-valued result will appear only on the first row of the result. The solution to this is to use the key field as part of a join against a simpler XPath query. As an example:

```sql
CREATE TABLE test (  
    id int PRIMARY KEY,  
    xml text  
);  
INSERT INTO test VALUES (1,  
    '<doc num="C1">  
    <line num="L1">\<a\>1</a><b\>2</b><c\>3</c></line>  
    <line num="L2">\<a\>11</a><b\>22</b><c\>33</c></line>  
  </doc>');  
INSERT INTO test VALUES (2,  
    '<doc num="C2">  
    <line num="L1">\<a\>111</a><b\>222</b><c\>333</c></line>  
    <line num="L2">\<a\>111</a><b\>222</b><c\>333</c></line>  
  </doc>');  
SELECT * FROM  
    xpath_table('id','xml','test',
    "xpath_string(article_xml,"/article/@date") > "2003-03-20" ")
    AS t(article_id integer, title text, author_id integer),
    tblPeopleInfo AS p
WHERE t.author_id = p.person_id;
```
Appendix F. Additional Supplied Modules

To get doc_num on every line, the solution is to use two invocations of xpath_table and join the results:

```sql
SELECT t.*, i.doc_num FROM
  xpath_table('id', 'xml', 'test', '
    /doc/line/@num|/doc/line/a|/doc/line/b|/doc/line/c', 'true')
  AS t(id int, line_num varchar(10), val1 int, val2 int, val3 int),
  xpath_table('id', 'xml', 'test', '/doc/@num', 'true')
  AS i(id int, doc_num varchar(10))
WHERE i.id=t.id AND i.id=1
ORDER BY doc_num, line_num;
```

<table>
<thead>
<tr>
<th>id</th>
<th>line_num</th>
<th>val1</th>
<th>val2</th>
<th>val3</th>
<th>doc_num</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>L1</td>
<td>1</td>
<td>2</td>
<td>3</td>
<td>C1</td>
</tr>
<tr>
<td>1</td>
<td>L2</td>
<td>11</td>
<td>22</td>
<td>33</td>
<td>C1</td>
</tr>
</tbody>
</table>

(2 rows)

F.46.4. XSLT Functions

The following functions are available if libxslt is installed:

F.46.4.1. xslt_process

```sql
xslt_process(text document, text stylesheet, text paramlist) returns text
```

This function applies the XSL stylesheet to the document and returns the transformed result. The paramlist is a list of parameter assignments to be used in the transformation, specified in the form a=1,b=2. Note that the parameter parsing is very simple-minded: parameter values cannot contain commas!

There is also a two-parameter version of xslt_process which does not pass any parameters to the transformation.
F.46.5. Author

John Gray <jgray@azuli.co.uk>

Development of this module was sponsored by Torchbox Ltd. (www.torchbox.com). It has the same BSD license as PostgreSQL.
Appendix G. Additional Supplied Programs

This appendix and the previous one contain information regarding the modules that can be found in the contrib directory of the PostgreSQL distribution. See Appendix F for more information about the contrib section in general and server extensions and plug-ins found in contrib specifically.

This appendix covers utility programs found in contrib. Once installed, either from source or a packaging system, they are found in the bin directory of the PostgreSQL installation and can be used like any other program.

G.1. Client Applications

This section covers PostgreSQL client applications in contrib. They can be run from anywhere, independent of where the database server resides. See also Reference II, PostgreSQL Client Applications for information about client applications that part of the core PostgreSQL distribution.

oid2name

Name

oid2name — resolve OIDs and file nodes in a PostgreSQL data directory

Synopsis

oid2name [option...]

Description

oid2name is a utility program that helps administrators to examine the file structure used by PostgreSQL. To make use of it, you need to be familiar with the database file structure, which is described in Chapter 65.

Note: The name "oid2name" is historical, and is actually rather misleading, since most of the time when you use it, you will really be concerned with tables’ filenode numbers (which are the file names visible in the database directories). Be sure you understand the difference between table OIDs and table filenodes!

oid2name connects to a target database and extracts OID, filenode, and/or table name information. You can also have it show database OIDs or tablespace OIDs.
Options

oid2name accepts the following command-line arguments:

- `f filenode`
  show info for table with filenode `filenode`
- `i`
  include indexes and sequences in the listing
- `o oid`
  show info for table with OID `oid`
- `q`
  omit headers (useful for scripting)
- `s`
  show tablespace OIDs
- `S`
  include system objects (those in `information_schema`, `pg_toast` and `pg_catalog` schemas)
- `t tablename_pattern`
  show info for table(s) matching ` tablename_pattern`
- `V` or `--version`
  Print the oid2name version and exit.
- `x`
  display more information about each object shown: tablespace name, schema name, and OID
- `?` or `--help`
  Show help about oid2name command line arguments, and exit.

oid2name also accepts the following command-line arguments for connection parameters:

- `d database`
  database to connect to
- `H host`
  database server’s host
- `p port`
  database server’s port
- `U username`
  user name to connect as
To display specific tables, select which tables to show by using \(-\text{o}, \text{-f}\) and/or \(-\text{t}\). \(-\text{o}\) takes an OID, \(-\text{f}\) takes a filenode, and \(-\text{t}\) takes a table name (actually, it’s a LIKE pattern, so you can use things like \text{foo\%}). You can use as many of these options as you like, and the listing will include all objects matched by any of the options. But note that these options can only show objects in the database given by \(-\text{d}\).

If you don’t give any of \(-\text{o}, \text{-f or -t}\), but do give \(-\text{d}\), it will list all tables in the database named by \(-\text{d}\). In this mode, the \(-\text{S}\) and \(-\text{i}\) options control what gets listed.

If you don’t give \(-\text{d}\) either, it will show a listing of database OIDs. Alternatively you can give \(-\text{s}\) to get a tablespace listing.

**Notes**

oid2name requires a running database server with non-corrupt system catalogs. It is therefore of only limited use for recovering from catastrophic database corruption situations.

**Examples**

```bash
$ # what’s in this database server, anyway?
$ oid2name
All databases:

<table>
<thead>
<tr>
<th>Oid</th>
<th>Database Name</th>
<th>Tablespace</th>
</tr>
</thead>
<tbody>
<tr>
<td>17228</td>
<td>alvherre</td>
<td>pg_default</td>
</tr>
<tr>
<td>17255</td>
<td>regression</td>
<td>pg_default</td>
</tr>
<tr>
<td>17227</td>
<td>template0</td>
<td>pg_default</td>
</tr>
<tr>
<td>1</td>
<td>template1</td>
<td>pg_default</td>
</tr>
</tbody>
</table>

$ oid2name -s
All tablespaces:

<table>
<thead>
<tr>
<th>Oid</th>
<th>Tablespace Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1663</td>
<td>pg_default</td>
</tr>
<tr>
<td>1664</td>
<td>pg_global</td>
</tr>
<tr>
<td>155151</td>
<td>fastdisk</td>
</tr>
<tr>
<td>155152</td>
<td>bigdisk</td>
</tr>
</tbody>
</table>

$ # OK, let’s look into database alvherre
$ cd SPGDATA/base/17228

$ # get top 10 db objects in the default tablespace, ordered by size
$ ls -ls * | head -10
```

-rw------- 1 alvherre alvherre 136536064 sep 14 09:51 155173
- rw------- 1 alvherre alvherre 17965056 sep 14 09:51 1155291
- rw------- 1 alvherre alvherre 1204224 sep 14 09:51 16717
- rw------- 1 alvherre alvherre 1204224 sep 14 09:51 16717
- rw------- 1 alvherre alvherre 581632 sep 6 17:51 1255
- rw------- 1 alvherre alvherre 237568 sep 14 09:50 16674
- rw------- 1 alvherre alvherre 212992 sep 14 09:51 1249
- rw------- 1 alvherre alvherre 204800 sep 14 09:51 16684

---

2515
$ # I wonder what file 155173 is ...
$ oid2name -d alvherre -f 155173
From database "alvherre":
<table>
<thead>
<tr>
<th>Filenode</th>
<th>Table Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>155173</td>
<td>accounts</td>
</tr>
</tbody>
</table>

$ # you can ask for more than one object
$ oid2name -d alvherre -f 155173 -f 1155291
From database "alvherre":
<table>
<thead>
<tr>
<th>Filenode</th>
<th>Table Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>155173</td>
<td>accounts</td>
</tr>
<tr>
<td>1155291</td>
<td>accounts_pkey</td>
</tr>
</tbody>
</table>

$ # you can mix the options, and get more details with -x
$ oid2name -d alvherre -t accounts -f 1155291 -x
From database "alvherre":
<table>
<thead>
<tr>
<th>Filenode</th>
<th>Table Name</th>
<th>Oid</th>
<th>Schema</th>
<th>Tablespace</th>
</tr>
</thead>
<tbody>
<tr>
<td>155173</td>
<td>accounts</td>
<td>155173</td>
<td>public</td>
<td>pg_default</td>
</tr>
<tr>
<td>1155291</td>
<td>accounts_pkey</td>
<td>1155291</td>
<td>public</td>
<td>pg_default</td>
</tr>
</tbody>
</table>

$ # show disk space for every db object
$ du [0-9]* | > while read SIZE FILENODE
> do
>  > echo "$$SIZE 'oid2name -q -d alvherre -i -f $FILENODE'"
> done
16          1155287  branches_pkey
16          1155289  tellers_pkey
17561       1155291  accounts_pkey
...

$ # same, but sort by size
$ du [0-9]* | sort -rn | while read SIZE FN
> do
>  > echo "$$SIZE 'oid2name -q -d alvherre -f $FN'"
> done
133466      155173  accounts
17561       1155291  accounts_pkey
1177        16717   pg_proc_priname_args_nsp_index
...

$ # If you want to see what’s in tablespaces, use the pg_tblspc directory
$ cd $PGDATA/pg_tblspc
$ oid2name -s
All tablespaces:
<table>
<thead>
<tr>
<th>Oid</th>
<th>Tablespace Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1663</td>
<td>pg_default</td>
</tr>
<tr>
<td>1664</td>
<td>pg_global</td>
</tr>
<tr>
<td>155151</td>
<td>fastdisk</td>
</tr>
</tbody>
</table>
155152 bigdisk

$ # what databases have objects in tablespace "fastdisk"?
$ ls -d 155151/
155151/17228/ 155151/PG_VERSION

$ # Oh, what was database 17228 again?
$ oid2name
All databases:

<table>
<thead>
<tr>
<th>Oid</th>
<th>Database Name</th>
<th>Tablespace</th>
</tr>
</thead>
<tbody>
<tr>
<td>17228</td>
<td>alvherre</td>
<td>pg_default</td>
</tr>
<tr>
<td>17255</td>
<td>regression</td>
<td>pg_default</td>
</tr>
<tr>
<td>17227</td>
<td>template0</td>
<td>pg_default</td>
</tr>
<tr>
<td>1</td>
<td>template1</td>
<td>pg_default</td>
</tr>
</tbody>
</table>

$ # Let’s see what objects does this database have in the tablespace.
$ cd 155151/17228
$ ls -l
total 0
-rw------- 1 postgres postgres 0 sep 13 23:20 155156

$ # OK, this is a pretty small table ... but which one is it?
$ oid2name -d alvherre -f 155156
From database "alvherre":

<table>
<thead>
<tr>
<th>Filenode</th>
<th>Table Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>155156</td>
<td>foo</td>
</tr>
</tbody>
</table>

Author

B. Palmer <bpalmer@crimelabs.net>
vacuumlo

Name
vacuumlo — remove orphaned large objects from a PostgreSQL database

Synopsis
vacuumlo [option...] dbname...

Description
vacuumlo is a simple utility program that will remove any "orphaned" large objects from a PostgreSQL database. An orphaned large object (LO) is considered to be any LO whose OID does not appear in any oid or lo data column of the database.

If you use this, you may also be interested in the lo_manage trigger in the lo module. lo_manage is useful to try to avoid creating orphaned LOs in the first place.

All databases named on the command line are processed.

Options
vacuumlo accepts the following command-line arguments:

- \( l \) limit
  Remove no more than \( limit \) large objects per transaction (default 1000). Since the server acquires a lock per LO removed, removing too many LOs in one transaction risks exceeding max_locks_per_transaction. Set the limit to zero if you want all removals done in a single transaction.

- n
  Don’t remove anything, just show what would be done.

- v
  Write a lot of progress messages.

- V
  --version
  Print the vacuumlo version and exit.

- ?
  --help
  Show help about vacuumlo command line arguments, and exit.

vacuumlo also accepts the following command-line arguments for connection parameters:
-h *hostname*
  Database server’s host.

-p *port*
  Database server’s port.

-U *username*
  User name to connect as.

-w
- --no-password
  Never issue a password prompt. If the server requires password authentication and a password is not available by other means such as a `.pgpass` file, the connection attempt will fail. This option can be useful in batch jobs and scripts where no user is present to enter a password.

-W
  Force vacuumlo to prompt for a password before connecting to a database.
  This option is never essential, since vacuumlo will automatically prompt for a password if the server demands password authentication. However, vacuumlo will waste a connection attempt finding out that the server wants a password. In some cases it is worth typing `-W` to avoid the extra connection attempt.

Notes

vacuumlo works by the following method: First, vacuumlo builds a temporary table which contains all of the OIDs of the large objects in the selected database. It then scans through all columns in the database that are of type `oid` or `lo`, and removes matching entries from the temporary table. (Note: Only types with these names are considered; in particular, domains over them are not considered.) The remaining entries in the temporary table identify orphaned LOs. These are removed.

Author

Peter Mount <peter@retep.org.uk>
G.2. Server Applications

This section covers PostgreSQL server-related applications in contrib. They are typically run on the host where the database server resides. See also Reference III, *PostgreSQL Server Applications* for information about server applications that part of the core PostgreSQL distribution.

**pg_standby**

**Name**

pg_standby — supports the creation of a PostgreSQL warm standby server

**Synopsis**

pg_standby [option...] archivelocation nextwalfile xlogfilepath [restartwalfile]

**Description**

pg_standby supports creation of a “warm standby” database server. It is designed to be a production-ready program, as well as a customizable template should you require specific modifications.

pg_standby is designed to be a waiting restore_command, which is needed to turn a standard archive recovery into a warm standby operation. Other configuration is required as well, all of which is described in the main server manual (see Section 26.2).

To configure a standby server to use pg_standby, put this into its recovery.conf configuration file:

```
restore_command = 'pg_standby archiveDir %f %p %r'
```

where `archiveDir` is the directory from which WAL segment files should be restored.

If `restartwalfile` is specified, normally by using the `%r` macro, then all WAL files logically preceding this file will be removed from `archivelocation`. This minimizes the number of files that need to be retained, while preserving crash-restart capability. Use of this parameter is appropriate if the `archivelocation` is a transient staging area for this particular standby server, but not when the `archivelocation` is intended as a long-term WAL archive area.

pg_standby assumes that `archivelocation` is a directory readable by the server-owning user. If `restartwalfile` (or `-k`) is specified, the `archivelocation` directory must be writable too.

There are two ways to fail over to a “warm standby” database server when the master server fails:

**Smart Failover**

In smart failover, the server is brought up after applying all WAL files available in the archive. This results in zero data loss, even if the standby server has fallen behind, but if there is a lot of unapplied WAL it can be a long time before the standby server becomes ready. To trigger a smart failover, create a trigger file containing the word `smart`, or just create it and leave it empty.

**Fast Failover**

In fast failover, the server is brought up immediately. Any WAL files in the archive that have not yet been applied will be ignored, and all transactions in those files are lost. To trigger a fast
failover, create a trigger file and write the word `fast` into it. `pg_standby` can also be configured to execute a fast failover automatically if no new WAL file appears within a defined interval.

**Options**

`pg_standby` accepts the following command-line arguments:

- `-c`
  
  Use `cp` or `copy` command to restore WAL files from archive. This is the only supported behavior so this option is useless.

- `-d`
  
  Print lots of debug logging output on `stderr`.

- `-k`
  
  Remove files from `archiveloaction` so that no more than this many WAL files before the current one are kept in the archive. Zero (the default) means not to remove any files from `archiveloaction`. This parameter will be silently ignored if `restartwalfile` is specified, since that specification method is more accurate in determining the correct archive cut-off point. Use of this parameter is `deprecated` as of PostgreSQL 8.3; it is safer and more efficient to specify a `restartwalfile` parameter. A too small setting could result in removal of files that are still needed for a restart of the standby server, while a too large setting wastes archive space.

- `-r maxretries`
  
  Set the maximum number of times to retry the copy command if it fails (default 3). After each failure, we wait for `sleeptime * num_retries` so that the wait time increases progressively. So by default, we will wait 5 secs, 10 secs, then 15 secs before reporting the failure back to the standby server. This will be interpreted as end of recovery and the standby will come up fully as a result.

- `-s sleeptime`
  
  Set the number of seconds (up to 60, default 5) to sleep between tests to see if the WAL file to be restored is available in the archive yet. The default setting is not necessarily recommended; consult Section 26.2 for discussion.

- `-t triggerfile`
  
  Specify a trigger file whose presence should cause failover. It is recommended that you use a structured file name to avoid confusion as to which server is being triggered when multiple servers exist on the same system; for example `/tmp/pgsql.trigger.5432`.

- `-V`

  Print the `pg_standby` version and exit.

- `-w maxwaittime`

  Set the maximum number of seconds to wait for the next WAL file, after which a fast failover will be performed. A setting of zero (the default) means wait forever. The default setting is not necessarily recommended; consult Section 26.2 for discussion.
-?
--help

Show help about pg_standby command line arguments, and exit.

Notes

pg_standby is designed to work with PostgreSQL 8.2 and later.

PostgreSQL 8.3 provides the %r macro, which is designed to let pg_standby know the last file it needs to keep. With PostgreSQL 8.2, the -k option must be used if archive cleanup is required. This option remains available in 8.3, but its use is deprecated.

PostgreSQL 8.4 provides the recovery_end_command option. Without this option a leftover trigger file can be hazardous.

pg_standby is written in C and has an easy-to-modify source code, with specifically designated sections to modify for your own needs.

Examples

On Linux or Unix systems, you might use:

archive_command = 'cp $p .../archive/$f'

restore_command = 'pg_standby -d -s 2 -t /tmp/pgsql.trigger.5442 .../archive %f $p %r 2>>standby.log'

recovery_end_command = 'rm -f /tmp/pgsql.trigger.5442'

where the archive directory is physically located on the standby server, so that the archive_command is accessing it across NFS, but the files are local to the standby (enabling use of ln). This will:

• produce debugging output in standby.log
• sleep for 2 seconds between checks for next WAL file availability
• stop waiting only when a trigger file called /tmp/pgsql.trigger.5442 appears, and perform failover according to its content
• remove the trigger file when recovery ends
• remove no-longer-needed files from the archive directory

On Windows, you might use:

archive_command = 'copy $p ...\archive\%f'

restore_command = 'pg_standby -d -s 5 -t C:\pgsql.trigger.5442 ...\archive %f $p %r 2>>standby.log'

recovery_end_command = 'del C:\pgsql.trigger.5442'

Note that backslashes need to be doubled in the archive_command, but not in the restore_command or recovery_end_command. This will:
• use the copy command to restore WAL files from archive
• produce debugging output in standby.log
• sleep for 5 seconds between checks for next WAL file availability
• stop waiting only when a trigger file called C:\pgsql.trigger.5442 appears, and perform failover according to its content
• remove the trigger file when recovery ends
• remove no-longer-needed files from the archive directory

The copy command on Windows sets the final file size before the file is completely copied, which would ordinarily confuse pg_standby. Therefore pg_standby waits sleeptime seconds once it sees the proper file size. GNUWin32’s cp sets the file size only after the file copy is complete.

Since the Windows example uses copy at both ends, either or both servers might be accessing the archive directory across the network.

Author
Simon Riggs <simon@2ndquadrant.com>

See Also
pg_archivecleanup
Appendix H. External Projects

PostgreSQL is a complex software project, and managing the project is difficult. We have found that many enhancements to PostgreSQL can be more efficiently developed separately from the core project.

H.1. Client Interfaces

There are only two client interfaces included in the base PostgreSQL distribution:

- libpq is included because it is the primary C language interface, and because many other client interfaces are built on top of it.
- ECPG is included because it depends on the server-side SQL grammar, and is therefore sensitive to changes in PostgreSQL itself.

All other language interfaces are external projects and are distributed separately. Table H-1 includes a list of some of these projects. Note that some of these packages might not be released under the same license as PostgreSQL. For more information on each language interface, including licensing terms, refer to its website and documentation.

Table H-1. Externally Maintained Client Interfaces

<table>
<thead>
<tr>
<th>Name</th>
<th>Language</th>
<th>Comments</th>
<th>Website</th>
</tr>
</thead>
<tbody>
<tr>
<td>DBD::Pg</td>
<td>Perl</td>
<td>Perl DBI driver</td>
<td><a href="https://metacpan.org/release/DBD-Pg/">https://metacpan.org/release/DBD-Pg/</a></td>
</tr>
<tr>
<td>JDBC</td>
<td>Java</td>
<td>Type 4 JDBC driver</td>
<td><a href="https://jdbc.postgresql.org/">https://jdbc.postgresql.org/</a></td>
</tr>
<tr>
<td>libpqxx</td>
<td>C++</td>
<td>C++ interface</td>
<td><a href="http://pqxx.org/">http://pqxx.org/</a></td>
</tr>
<tr>
<td>node-postgres</td>
<td>JavaScript</td>
<td>Node.js driver</td>
<td><a href="https://node-postgres.com/">https://node-postgres.com/</a></td>
</tr>
<tr>
<td>Npgsql</td>
<td>.NET</td>
<td>.NET data provider</td>
<td><a href="http://www.npgsql.org/">http://www.npgsql.org/</a></td>
</tr>
<tr>
<td>pgtcl</td>
<td>Tcl</td>
<td></td>
<td><a href="https://github.com/flightaware/Pgtcl">https://github.com/flightaware/Pgtcl</a></td>
</tr>
<tr>
<td>pgtclng</td>
<td>Tcl</td>
<td></td>
<td><a href="http://sourceforge.net/projects/pgtclng/">http://sourceforge.net/projects/pgtclng/</a></td>
</tr>
<tr>
<td>pq</td>
<td>Go</td>
<td>Pure Go driver for Go’s database/sql</td>
<td><a href="https://github.com/libpq">https://github.com/libpq</a></td>
</tr>
<tr>
<td>psycopg</td>
<td>Python</td>
<td>DB API 2.0-compliant</td>
<td><a href="http://initd.org/psycopg/">http://initd.org/psycopg/</a></td>
</tr>
<tr>
<td>psqlODBC</td>
<td>ODBC</td>
<td>ODBC driver</td>
<td><a href="https://odbc.postgresql.org/">https://odbc.postgresql.org/</a></td>
</tr>
</tbody>
</table>
Appendix H. External Projects

H.2. Administration Tools

There are several administration tools available for PostgreSQL. The most popular is pgAdmin\(^1\), and there are several commercially available ones as well.

H.3. Procedural Languages

PostgreSQL includes several procedural languages with the base distribution: PL/pgSQL, PL/Tcl, PL/Perl, and PL/Python.

In addition, there are a number of procedural languages that are developed and maintained outside the core PostgreSQL distribution. Table H-2 lists some of these packages. Note that some of these projects might not be released under the same license as PostgreSQL. For more information on each procedural language, including licensing information, refer to its website and documentation.

<table>
<thead>
<tr>
<th>Name</th>
<th>Language</th>
<th>Website</th>
</tr>
</thead>
<tbody>
<tr>
<td>PL/Java</td>
<td>Java</td>
<td><a href="https://github.com/tada/pljava">https://github.com/tada/pljava</a></td>
</tr>
<tr>
<td>PL/PHP</td>
<td>PHP</td>
<td><a href="https://public.commandprompt.com/projects/plphp">https://public.commandprompt.com/projects/plphp</a></td>
</tr>
<tr>
<td>PL/Py</td>
<td>Python</td>
<td><a href="http://python.projects.postgresql.org/backend/">http://python.projects.postgresql.org/backend/</a></td>
</tr>
<tr>
<td>PL/R</td>
<td>R</td>
<td><a href="https://github.com/postgres-plr/plr">https://github.com/postgres-plr/plr</a></td>
</tr>
<tr>
<td>PL/Ruby</td>
<td>Ruby</td>
<td><a href="http://raa.ruby-lang.org/project/pl-ruby/">http://raa.ruby-lang.org/project/pl-ruby/</a></td>
</tr>
<tr>
<td>PL/Scheme</td>
<td>Scheme</td>
<td><a href="http://plescheme.projects.postgresql.org/">http://plescheme.projects.postgresql.org/</a></td>
</tr>
<tr>
<td>PL/sh</td>
<td>Unix shell</td>
<td><a href="https://github.com/petere/plsh">https://github.com/petere/plsh</a></td>
</tr>
</tbody>
</table>

H.4. Extensions

PostgreSQL is designed to be easily extensible. For this reason, extensions loaded into the database can function just like features that are built in. The contrib/ directory shipped with the source code contains several extensions, which are described in Appendix F. Other extensions are developed independently, like PostGIS\(^2\). Even PostgreSQL replication solutions can be developed externally. For example, Slony-I\(^3\) is a popular master/standby replication solution that is developed independently from the core project.

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2. http://postgis.net/
Appendix I. The Source Code Repository

The PostgreSQL source code is stored and managed using the Git version control system. A public mirror of the master repository is available; it is updated within a minute of any change to the master repository.

Our wiki, https://wiki.postgresql.org/wiki/Working_with_Git, has some discussion on working with Git.

Note that building PostgreSQL from the source repository requires reasonably up-to-date versions of bison, flex, and Perl. These tools are not needed to build from a distribution tarball, because the files that these tools are used to build are included in the tarball. Other tool requirements are the same as shown in Section 16.2.

I.1. Getting The Source via Git

With Git you will make a copy of the entire code repository on your local machine, so you will have access to all history and branches offline. This is the fastest and most flexible way to develop or test patches.

Git

1. You will need an installed version of Git, which you can get from https://git-scm.com. Many systems already have a recent version of Git installed by default, or available in their package distribution system.

2. To begin using the Git repository, make a clone of the official mirror:

   
   git clone https://git.postgresql.org/git/postgresql.git
   
   This will copy the full repository to your local machine, so it may take a while to complete, especially if you have a slow Internet connection. The files will be placed in a new subdirectory postgresql of your current directory.

   The Git mirror can also be reached via the Git protocol. Just change the URL prefix to git, as in:

   
   git clone git://git.postgresql.org/git/postgresql.git
   
3. Whenever you want to get the latest updates in the system, cd into the repository, and run:

   git fetch

Git can do a lot more things than just fetch the source. For more information, consult the Git man pages, or see the website at https://git-scm.com.
Appendix J. Documentation

PostgreSQL has four primary documentation formats:

- Plain text, for pre-installation information
- HTML, for on-line browsing and reference
- PDF or PostScript, for printing
- man pages, for quick reference.

Additionally, a number of plain-text README files can be found throughout the PostgreSQL source tree, documenting various implementation issues.

HTML documentation and man pages are part of a standard distribution and are installed by default. PDF and PostScript format documentation is available separately for download.

J.1. DocBook

The documentation sources are written in DocBook, which is a markup language superficially similar to HTML. Both of these languages are applications of the Standard Generalized Markup Language, SGML, which is essentially a language for describing other languages. In what follows, the terms DocBook and SGML are both used, but technically they are not interchangeable.

DocBook allows an author to specify the structure and content of a technical document without worrying about presentation details. A document style defines how that content is rendered into one of several final forms. DocBook is maintained by the OASIS group\(^1\). The official DocBook site\(^2\) has good introductory and reference documentation and a complete O’Reilly book for your online reading pleasure. The NewbieDoc Docbook Guide\(^3\) is very helpful for beginners. The FreeBSD Documentation Project\(^4\) also uses DocBook and has some good information, including a number of style guidelines that might be worth considering.

J.2. Tool Sets

The following tools are used to process the documentation. Some might be optional, as noted.

DocBook DTD\(^5\)

This is the definition of DocBook itself. We currently use version 4.2; you cannot use later or earlier versions. You need the SGML variant of the DocBook DTD, but to build man pages you also need the XML variant of the same version.

---

ISO 8879 character entities\(^6\)

These are required by DocBook but are distributed separately because they are maintained by ISO.

DocBook DSSSL Stylesheets\(^7\)

These contain the processing instructions for converting the DocBook sources to other formats, such as HTML.

DocBook XSL Stylesheets\(^8\)

This is another stylesheet for converting DocBook to other formats. We currently use this to produce man pages and optionally HTMLHelp. You can also use this toolchain to produce HTML or PDF output, but official PostgreSQL releases use the DSSSL stylesheets for that.

The minimum required version is currently 1.74.0.

OpenJade\(^9\)

This is the base package of SGML processing. It contains an SGML parser, a DSSSL processor (that is, a program to convert SGML to other formats using DSSSL stylesheets), as well as a number of related tools. Jade is now being maintained by the OpenJade group, no longer by James Clark.

Libxml2\(^10\) for `xmllint`

This library and the `xmllint` tool it contains are used for processing XML. Many developers will already have Libxml2 installed, because it is also used when building the PostgreSQL code. Note, however, that `xmllint` might need to be installed from a separate subpackage.

Libxslt\(^11\) for `xsltproc`

This is the processing tool to use with the XSLT stylesheets (like jade is the processing tool for DSSSL stylesheets).

JadeTeX\(^12\)

If you want to, you can also install JadeTeX to use TeX as a formatting backend for Jade. JadeTeX can create PostScript or PDF files (the latter with bookmarks).

However, the output from JadeTeX is inferior to what you get from the RTF backend. Particular problem areas are tables and various artifacts of vertical and horizontal spacing. Also, there is no opportunity to manually polish the results.

We have documented experience with several installation methods for the various tools that are needed to process the documentation. These will be described below. There might be some other packaged distributions for these tools. Please report package status to the documentation mailing list, and we will include that information here.

\[^6\] http://www.oasis-open.org/cover/ISOEnts.zip
\[^7\] http://wiki.docbook.org/DocBookDssslStylesheetDocs
\[^8\] http://wiki.docbook.org/DocBookXslStylesheets
\[^9\] http://openjade.sourceforge.net
\[^10\] http://xmlsoft.org/
\[^11\] http://xmlsoft.org/XSLT/
\[^12\] http://jadetex.sourceforge.net
Appendix J. Documentation

J.2.1. Installation on Fedora, RHEL, and Derivatives

To install the required packages, use:

`yum install docbook-dtds docbook-style-dsssl docbook-style-xsl libxslt openjade`

J.2.2. Installation on FreeBSD

The FreeBSD Documentation Project is itself a heavy user of DocBook, so it comes as no surprise that there is a full set of “ports” of the documentation tools available on FreeBSD. The following ports need to be installed to build the documentation on FreeBSD.

- `textproc/docbook-sgml`
- `textproc/docbook-xml`
- `textproc/docbook-xsl`
- `textproc/dsssl-docbook-modular`
- `textproc/libxslt`
- `textproc/openjade`

A number of things from `/usr/ports/print (tex, jadetex)` might also be of interest.

More information about the FreeBSD documentation tools can be found in the FreeBSD Documentation Project’s instructions\(^{13}\).

J.2.3. Debian Packages

There is a full set of packages of the documentation tools available for Debian GNU/Linux. To install, simply use:

`apt-get install docbook docbook-dsssl docbook-xsl libxml2-utils openjade1.3 opensp xsltproc`

J.2.4. OS X

If you use MacPorts, the following will get you set up:

`sudo port install docbook-dsssl docbook-sgml-4.2 docbook-xml-4.2 docbook-xsl libxslt opensp xsltproc`

Appendix J. Documentation

J.2.5. Manual Installation from Source

The manual installation process of the DocBook tools is somewhat complex, so if you have pre-built packages available, use them. We describe here only a standard setup, with reasonably standard installation paths, and no “fancy” features. For details, you should study the documentation of the respective package, and read SGML introductory material.

J.2.5.1. Installing OpenJade

1. The installation of OpenJade offers a GNU-style ./configure; make; make install build process. Details can be found in the OpenJade source distribution. In a nutshell:

   ./configure --enable-default-catalog=/usr/local/share/sgml/catalog
   make
   make install

   Be sure to remember where you put the “default catalog”; you will need it below. You can also leave it off, but then you will have to set the environment variable SGML_CATALOG_FILES to point to the file whenever you use jade later on. (This method is also an option if OpenJade is already installed and you want to install the rest of the toolchain locally.)

   **Note:** Some users have reported encountering a segmentation fault using OpenJade 1.4development to build the PDFs, with a message like:

   openjade:./stylesheet.dsl:664:2:E: flow object not accepted by port; only display flow objects
   Downgrading to OpenJade 1.3 should get rid of this error.

2. Additionally, you should install the files dsssl.dtd, fot.dtd, style-sheet.dtd, and catalog from the dsssl directory somewhere, perhaps into /usr/local/share/sgml/dsssl. It’s probably easiest to copy the entire directory:

   cp -R dsssl /usr/local/share/sgml

3. Finally, create the file /usr/local/share/sgml/catalog and add this line to it:

   CATALOG "dsssl/catalog"

   (This is a relative path reference to the file installed in step 2. Be sure to adjust it if you chose your installation layout differently.)

J.2.5.2. Installing the DocBook DTD Kit

1. Obtain the DocBook V4.2 distribution\(^\text{14}\).

2. Create the directory /usr/local/share/sgml/docbook-4.2 and change to it. (The exact location is irrelevant, but this one is reasonable within the layout we are following here.)

   $ mkdir /usr/local/share/sgml/docbook-4.2
   $ cd /usr/local/share/sgml/docbook-4.2

3. Unpack the archive:

   $ unzip -a ...../docbook-4.2.zip

   (The archive will unpack its files into the current directory.)

\(^\text{14}\) http://www.docbook.org/sgml/4.2/docbook-4.2.zip
4. Edit the file /usr/local/share/sgml/catalog (or whatever you told jade during installation) and put a line like this into it:
   
   CATALOG "docbook-4.2/docbook.cat"

5. Download the ISO 8879 character entities archive, unpack it, and put the files in the same directory you put the DocBook files in:
   
   $ cd /usr/local/share/sgml/docbook-4.2
   $ unzip ...../ISOEnts.zip

6. Run the following command in the directory with the DocBook and ISO files:
   
   perl -pi -e 's/iso-\(.*\).gml/ISO\1/g' docbook.cat
   
   (This fixes a mixup between the names used in the DocBook catalog file and the actual names of the ISO character entity files.)

J.2.5.3. Installing the DocBook DSSSL Style Sheets

To install the style sheets, unzip and untar the distribution and move it to a suitable place, for example /usr/local/share/sgml. (The archive will automatically create a subdirectory.)

   $ gunzip docbook-dsssl-1.xx.tar.gz
   $ tar -C /usr/local/share/sgml -xf docbook-dsssl-1.xx.tar

The usual catalog entry in /usr/local/share/sgml/catalog can also be made:

   CATALOG "docbook-dsssl-1.xx/catalog"

Because stylesheets change rather often, and it’s sometimes beneficial to try out alternative versions, PostgreSQL doesn’t use this catalog entry. See Section J.2.6 for information about how to select the stylesheets instead.

J.2.5.4. Installing JadeTeX

To install and use JadeTeX, you will need a working installation of TeX and LaTeX2e, including the supported tools and graphics packages, Babel, AMS fonts and AMS-LaTeX, the PSNFSS extension and companion kit of “the 35 fonts”, the dvips program for generating PostScript, the macro packages fancyhdr, hyperref, minitoc, url and ot2enc. All of these can be found on your friendly neighborhood CTAN site. The installation of the TeX base system is far beyond the scope of this introduction. Binary packages should be available for any system that can run TeX.

Before you can use JadeTeX with the PostgreSQL documentation sources, you will need to increase the size of TeX’s internal data structures. Details on this can be found in the JadeTeX installation instructions.

Once that is finished you can install JadeTeX:

   $ gunzip jadetex-xxx.tar.gz
   $ tar xf jadetex-xxx.tar
   $ cd jadetex
   $ make install
   $ mktexlsr

Appendix J. Documentation

The last two need to be done as root.

J.2.6. Detection by configure

Before you can build the documentation you need to run the configure script as you would when building the PostgreSQL programs themselves. Check the output near the end of the run, it should look something like this:

```
checking for onsgmls... onsgmls
checking for openjade... openjade
checking for DocBook V4.2... yes
checking for DocBook stylesheets... /usr/share/sgml/docbook/stylesheet/dsssl/modular
checking for collateindex.pl... /usr/bin/collateindex.pl
checking for xsltproc... xsltproc
checking for osx... osx
```

If neither `onsgmls` nor `nsgmls` were found then some of the following tests will be skipped. `nsgmls` is part of the Jade package. You can pass the environment variables `JADE` and `NSGMLS` to configure to point to the programs if they are not found automatically. If “DocBook V4.2” was not found then you did not install the DocBook DTD kit in a place where Jade can find it, or you have not set up the catalog files correctly. See the installation hints above. The DocBook stylesheets are looked for in a number of relatively standard places, but if you have them some other place then you should set the environment variable `DOCBOOKSTYLE` to the location and rerun configure afterwards.

J.3. Building The Documentation

Once you have everything set up, change to the directory `doc/src/sgml` and run one of the commands described in the following subsections to build the documentation. (Remember to use GNU make.)

J.3.1. HTML

To build the HTML version of the documentation:

```
doc/src/sgml$ make html
```

This is also the default target. The output appears in the subdirectory `html`.

To produce HTML documentation with the stylesheet used on postgresql.org\(^\text{17}\) instead of the default simple style use:

```
doc/src/sgml$ make STYLE=website html
```

To create a proper index, the build might process several identical stages. If you do not care about the index, and just want to proof-read the output, use `draft`:

```
doc/src/sgml$ make draft
```

\(^\text{17}\) https://www.postgresql.org/docs/current
To build the documentation as a single HTML page, use:

```
doc/src/sgml$ make postgres.html
```

### J.3.2. Manpages

We use the DocBook XSL stylesheets to convert DocBook `refentry` pages to *roff output suitable for man pages. The man pages are also distributed as a tar archive, similar to the HTML version. To create the man pages, use the commands:

```
cd doc/src/sgml
make man
```

### J.3.3. Print Output via JadeTeX

If you want to use JadeTeX to produce a printable rendition of the documentation, you can use one of the following commands:

- To generate PostScript via DVI in A4 format:
  ```
doc/src/sgml$ make postgres-A4.ps
```
- In U.S. letter format:
  ```
doc/src/sgml$ make postgres-US.ps
```
- To make a PDF:
  ```
doc/src/sgml$ make postgres-A4.pdf
```
  or:
  ```
doc/src/sgml$ make postgres-US.pdf
```
  (Of course you can also make a PDF version from the PostScript, but if you generate PDF directly, it will have hyperlinks and other enhanced features.)

When using JadeTeX to build the PostgreSQL documentation, you will probably need to increase some of TeX’s internal parameters. These can be set in the file `texmf.cnf`. The following settings worked at the time of this writing:

```
hash_extra.jadetex = 200000
hash_extra.pdfjadetex = 200000
pool_size.jadetex = 2000000
pool_size.pdfjadetex = 2000000
string_vacancies.jadetex = 150000
string_vacancies.pdfjadetex = 150000
max_strings.jadetex = 300000
max_strings.pdfjadetex = 300000
save_size.jadetex = 15000
save_size.pdfjadetex = 15000
```
Appendix J. Documentation

J.3.4. Overflow Text

Occasionally text is too wide for the printed margins, and in extreme cases, too wide for the printed page, e.g. non-wrapped text, wide tables. Overly wide text generates “Overfull hbox” messages in the TeX log output file, e.g. postgres-US.log or postgres-A4.log. There are 72 points in an inch so anything reported as over 72 points too wide will probably not fit on the printed page (assuming one inch margins). To find the SGML text causing the overflow, find the first page number mentioned above the overflow message, e.g. [50 ###] (page 50), and look at the page after that (e.g. page 51) in the PDF file to see the overflow text and adjust the SGML accordingly.

J.3.5. Print Output via RTF

You can also create a printable version of the PostgreSQL documentation by converting it to RTF and applying minor formatting corrections using an office suite. Depending on the capabilities of the particular office suite, you can then convert the documentation to PostScript of PDF. The procedure below illustrates this process using Applixware.

Note: It appears that current versions of the PostgreSQL documentation trigger some bug in or exceed the size limit of OpenJade. If the build process of the RTF version hangs for a long time and the output file still has size 0, then you might have hit that problem. (But keep in mind that a normal build takes 5 to 10 minutes, so don’t abort too soon.)

Applixware RTF Cleanup

OpenJade omits specifying a default style for body text. In the past, this undiagnosed problem led to a long process of table of contents generation. However, with great help from the Applixware folks the symptom was diagnosed and a workaround is available.

1. Generate the RTF version by typing:
   
   doc/src/sgml$ make postgres.rtf

2. Repair the RTF file to correctly specify all styles, in particular the default style. If the document contains refentry sections, one must also replace formatting hints which tie a preceding paragraph to the current paragraph, and instead tie the current paragraph to the following one. A utility, fixrtf, is available in doc/src/sgml to accomplish these repairs:
   
   doc/src/sgml$ ./fixrtf --refentry postgres.rtf

   The script adds \{s0 Normal;\} as the zeroth style in the document. According to Applixware, the RTF standard would prohibit adding an implicit zeroth style, though Microsoft Word happens to handle this case. For repairing refentry sections, the script replaces \keepn tags with \keep.

3. Open a new document in Applixware Words and then import the RTF file.

4. Generate a new table of contents (ToC) using Applixware.
   
   a. Select the existing ToC lines, from the beginning of the first character on the first line to the last character of the last line.
Appendix J. Documentation

b. Build a new ToC using Tools → Book Building → Create Table of Contents. Select the first three levels of headers for inclusion in the ToC. This will replace the existing lines imported in the RTF with a native Applixware ToC.

c. Adjust the ToC formatting by using Format → Style, selecting each of the three ToC styles, and adjusting the indents for First and Left. Use the following values:

<table>
<thead>
<tr>
<th>Style</th>
<th>First Indent (inches)</th>
<th>Left Indent (inches)</th>
</tr>
</thead>
<tbody>
<tr>
<td>TOC-Heading 1</td>
<td>0.4</td>
<td>0.4</td>
</tr>
<tr>
<td>TOC-Heading 2</td>
<td>0.8</td>
<td>0.8</td>
</tr>
<tr>
<td>TOC-Heading 3</td>
<td>1.2</td>
<td>1.2</td>
</tr>
</tbody>
</table>

5. Work through the document to:

- Adjust page breaks.
- Adjust table column widths.

6. Replace the right-justified page numbers in the Examples and Figures portions of the ToC with correct values. This only takes a few minutes.

7. Delete the index section from the document if it is empty.

8. Regenerate and adjust the table of contents.

   a. Select the ToC field.
   b. Select Tools → Book Building → Create Table of Contents.
   c. Unbind the ToC by selecting Tools → Field Editing → Unprotect.
   d. Delete the first line in the ToC, which is an entry for the ToC itself.

9. Save the document as native Applixware Words format to allow easier last minute editing later.

10. “Print” the document to a file in PostScript format.

J.3.6. Plain Text Files

The installation instructions are also distributed as plain text, in case they are needed in a situation where better reading tools are not available. The INSTALL file corresponds to Chapter 16, with some minor changes to account for the different context. To recreate the file, change to the directory doc/src/sgml and enter `make INSTALL`.

In the past, the release notes and regression testing instructions were also distributed as plain text, but this practice has been discontinued.

J.3.7. Syntax Check

Building the documentation can take very long. But there is a method to just check the correct syntax of the documentation files, which only takes a few seconds:

doc/src/sgml$ make check
J.4. Documentation Authoring

SGML and DocBook do not suffer from an oversupply of open-source authoring tools. The most common tool set is the Emacs/XEmacs editor with appropriate editing mode. On some systems these tools are provided in a typical full installation.

J.4.1. Emacs/PSGML

PSGML is the most common and most powerful mode for editing SGML documents. When properly configured, it will allow you to use Emacs to insert tags and check markup consistency. You could use it for HTML as well. Check the PSGML web site\textsuperscript{18} for downloads, installation instructions, and detailed documentation.

There is one important thing to note with PSGML: its author assumed that your main SGML DTD directory would be "/usr/local/lib/sgml". If, as in the examples in this chapter, you use "/usr/local/share/sgml", you have to compensate for this, either by setting \texttt{SGML\_CATALOG\_FILES} environment variable, or you can customize your PSGML installation (its manual tells you how).

Put the following in your "/.emacs" environment file (adjusting the path names to be appropriate for your system):

```lisp
; ********** for SGML mode (psgml)
(setq sgml-omittag t)
(setq sgml-shorttag t)
(setq sgml-minimize-attributes nil)
(setq sgml-always-quote-attributes t)
(setq sgml-indent-step 1)
(setq sgml-indent-data t)
(setq sgml-parent-document nil)
(setq sgml-exposed-tags nil)
(setq sgml-catalog-files 
"(/usr/local/share/sgml/catalog")

(autoload 'sgml-mode "psgml" "Major mode to edit SGML files." t )
```

and in the same file add an entry for SGML into the (existing) definition for \texttt{auto-mode-alist}:

```lisp
(setq
  auto-mode-alist
  '(("\.sgml" . sgml-mode)
    ))
```

You might find that when using PSGML, a comfortable way of working with these separate files of book parts is to insert a proper \texttt{DOCTYPE} declaration while you’re editing them. If you are working on this source, for instance, it is an appendix chapter, so you would specify the document as an “appendix” instance of a DocBook document by making the first line look like this:

\footnote{http://www.lysator.liu.se/projects/about_psgml.html}
This means that anything and everything that reads SGML will get it right, and I can verify the document with `nsgmls -s docguide.sgml`. (But you need to take out that line before building the entire documentation set.)

### J.4.2. Other Emacs Modes

GNU Emacs ships with a different SGML mode, which is not quite as powerful as PSGML, but it’s less confusing and lighter weight. Also, it offers syntax highlighting (font lock), which can be very helpful. `src/tools/editors/emacs.samples` contains sample settings for this mode.

Norm Walsh offers a major mode specifically for DocBook which also has font-lock and a number of features to reduce typing.

### J.5. Style Guide

#### J.5.1. Reference Pages

Reference pages should follow a standard layout. This allows users to find the desired information more quickly, and it also encourages writers to document all relevant aspects of a command. Consistency is not only desired among PostgreSQL reference pages, but also with reference pages provided by the operating system and other packages. Hence the following guidelines have been developed. They are for the most part consistent with similar guidelines established by various operating systems.

Reference pages that describe executable commands should contain the following sections, in this order. Sections that do not apply can be omitted. Additional top-level sections should only be used in special circumstances; often that information belongs in the “Usage” section.

**Name**

This section is generated automatically. It contains the command name and a half-sentence summary of its functionality.

**Synopsis**

This section contains the syntax diagram of the command. The synopsis should normally not list each command-line option; that is done below. Instead, list the major components of the command line, such as where input and output files go.

**Description**

Several paragraphs explaining what the command does.

**Options**

A list describing each command-line option. If there are a lot of options, subsections can be used.

**Exit Status**

If the program uses 0 for success and non-zero for failure, then you do not need to document it. If there is a meaning behind the different non-zero exit codes, list them here.

Appendix J. Documentation

Usage
Describe any sublanguage or run-time interface of the program. If the program is not interactive, this section can usually be omitted. Otherwise, this section is a catch-all for describing run-time features. Use subsections if appropriate.

Environment
List all environment variables that the program might use. Try to be complete; even seemingly trivial variables like SHELL might be of interest to the user.

Files
List any files that the program might access implicitly. That is, do not list input and output files that were specified on the command line, but list configuration files, etc.

Diagnostics
Explain any unusual output that the program might create. Refrain from listing every possible error message. This is a lot of work and has little use in practice. But if, say, the error messages have a standard format that the user can parse, this would be the place to explain it.

Notes
Anything that doesn’t fit elsewhere, but in particular bugs, implementation flaws, security considerations, compatibility issues.

Examples
Examples

History
If there were some major milestones in the history of the program, they might be listed here. Usually, this section can be omitted.

Author
Author (only used in the contrib section)

See Also
Cross-references, listed in the following order: other PostgreSQL command reference pages, PostgreSQL SQL command reference pages, citation of PostgreSQL manuals, other reference pages (e.g., operating system, other packages), other documentation. Items in the same group are listed alphabetically.

Reference pages describing SQL commands should contain the following sections: Name, Synopsis, Description, Parameters, Outputs, Notes, Examples, Compatibility, History, See Also. The Parameters section is like the Options section, but there is more freedom about which clauses of the command can be listed. The Outputs section is only needed if the command returns something other than a default command-completion tag. The Compatibility section should explain to what extent this command conforms to the SQL standard(s), or to which other database system it is compatible. The See Also section of SQL commands should list SQL commands before cross-references to programs.
Appendix K. Acronyms

This is a list of acronyms commonly used in the PostgreSQL documentation and in discussions about PostgreSQL.

ANSI
American National Standards Institute\(^1\)

API
Application Programming Interface\(^2\)

ASCII
American Standard Code for Information Interchange\(^3\)

BKI
Backend Interface

CA
Certificate Authority\(^4\)

CIDR
Classless Inter-Domain Routing\(^5\)

CPAN
Comprehensive Perl Archive Network\(^6\)

CRL
Certificate Revocation List\(^7\)

CSV
Comma Separated Values\(^8\)

CTE
Common Table Expression

CVE
Common Vulnerabilities and Exposures\(^9\)

DBA
Database Administrator\(^{10}\)

6. [https://www.cpan.org/](https://www.cpan.org/)
Appendix K. Acronyms

DBI
Database Interface (Perl)\textsuperscript{11}

DBMS
Database Management System\textsuperscript{12}

DDL
Data Definition Language\textsuperscript{13}, SQL commands such as \texttt{CREATE TABLE}, \texttt{ALTER USER}

DML
Data Manipulation Language\textsuperscript{14}, SQL commands such as \texttt{INSERT}, \texttt{UPDATE}, \texttt{DELETE}

DST
Daylight Saving Time\textsuperscript{15}

ECPG
Embedded C for PostgreSQL

ESQL
Embedded SQL\textsuperscript{16}

FAQ
Frequently Asked Questions\textsuperscript{17}

FSM
Free Space Map

GEQO
Genetic Query Optimizer

GIN
Generalized Inverted Index

GiST
Generalized Search Tree

Git
Git\textsuperscript{18}

GMT
Greenwich Mean Time\textsuperscript{19}

GSSAPI
Generic Security Services Application Programming Interface\textsuperscript{20}

GUC
Grand Unified Configuration, the PostgreSQL subsystem that handles server configuration

\textsuperscript{11} http://dbi.perl.org/
\textsuperscript{12} http://en.wikipedia.org/wiki/Dbms
\textsuperscript{13} http://en.wikipedia.org/wiki/Data_Definition_Language
\textsuperscript{14} http://en.wikipedia.org/wiki/Data_Manipulation_Language
\textsuperscript{15} http://en.wikipedia.org/wiki/Daylight_saving_time
\textsuperscript{16} http://en.wikipedia.org/wiki/Embedded_SQL
\textsuperscript{17} http://en.wikipedia.org/wiki/FAQ
\textsuperscript{18} http://en.wikipedia.org/wiki/Git\,(software)
\textsuperscript{19} http://en.wikipedia.org/wiki/GMT
\textsuperscript{20} http://en.wikipedia.org/wiki/Generic_Security_Services_Application_Program_Interface
Appendix K. Acronyms

HBA
Host-Based Authentication

HOT
Heap-Only Tuples<sup>21</sup>

IEC
International Electrotechnical Commission<sup>22</sup>

IEEE
Institute of Electrical and Electronics Engineers<sup>23</sup>

IPC
Inter-Process Communication<sup>24</sup>

ISO
International Organization for Standardization<sup>25</sup>

ISSN
International Standard Serial Number<sup>26</sup>

JDBC
Java Database Connectivity<sup>27</sup>

LDAP
Lightweight Directory Access Protocol<sup>28</sup>

MSVC
Microsoft Visual C<sup>29</sup>

MVCC
Multi-Version Concurrency Control

NLS
National Language Support<sup>30</sup>

ODBC
Open Database Connectivity<sup>31</sup>

OID
Object Identifier

OLAP
Online Analytical Processing<sup>32</sup>

---

<sup>21</sup> https://git.postgresql.org/gitweb?p=postgresql.git;a=blob;f=src/backend/access/heap/README.HOT;hb=HEAD
<sup>22</sup> http://en.wikipedia.org/wiki/International_Electrotechnical_Commission
<sup>23</sup> http://standards.ieee.org/
<sup>24</sup> http://en.wikipedia.org/wiki/Inter-process_communication
<sup>25</sup> http://www.iso.org/iso/home.htm
<sup>26</sup> http://en.wikipedia.org/wiki/Issn
<sup>27</sup> http://en.wikipedia.org/wiki/Java_Database_Connectivity
<sup>29</sup> http://en.wikipedia.org/wiki/Visual_C++
<sup>30</sup> http://en.wikipedia.org/wiki/Internationalization_and_localization
<sup>31</sup> http://en.wikipedia.org/wiki/Open_Database_Connectivity
<sup>32</sup> http://en.wikipedia.org/wiki/Olap
Appendix K. Acronyms

OLTP
Online Transaction Processing\(^\text{33}\)

ORDBMS
Object-Relational Database Management System\(^\text{34}\)

PAM
Pluggable Authentication Modules\(^\text{35}\)

PGSQL
PostgreSQL

PGXS
PostgreSQL Extension System

PID
Process Identifier\(^\text{36}\)

PITR
Point-In-Time Recovery (Continuous Archiving)

PL
Procedural Languages (server-side)

POSIX
Portable Operating System Interface\(^\text{37}\)

RDBMS
Relational Database Management System\(^\text{38}\)

RFC
Request For Comments\(^\text{39}\)

SGML
Standard Generalized Markup Language\(^\text{40}\)

SPI
Server Programming Interface

SP-GiST
Space-Partitioned Generalized Search Tree

SQL
Structured Query Language\(^\text{41}\)

SRF
Set-Returning Function

\(^{33}\) http://en.wikipedia.org/wiki/OLTP
\(^{34}\) http://en.wikipedia.org/wiki/ORDBMS
\(^{35}\) http://en.wikipedia.org/wiki/Pluggable_Authentication_Modules
\(^{36}\) http://en.wikipedia.org/wiki/Process_identifier
\(^{37}\) http://en.wikipedia.org/wiki/POSIX
\(^{38}\) http://en.wikipedia.org/wiki/Relational_database_management_system
\(^{39}\) http://en.wikipedia.org/wiki/Request_for_Comments
\(^{40}\) http://en.wikipedia.org/wiki/SGML
\(^{41}\) http://en.wikipedia.org/wiki/SQL
Appendix K. Acronyms

SSH
  Secure Shell\(^{42}\)

SSL
  Secure Sockets Layer\(^{43}\)

SSPI
  Security Support Provider Interface\(^{44}\)

SYSV
  Unix System V\(^{45}\)

TCP/IP
  Transmission Control Protocol (TCP) / Internet Protocol (IP)\(^{46}\)

TID
  Tuple Identifier

TOAST
  The Oversized-Attribute Storage Technique

TPC
  Transaction Processing Performance Council\(^{47}\)

URL
  Uniform Resource Locator\(^{48}\)

UTC
  Coordinated Universal Time\(^{49}\)

UTF
  Unicode Transformation Format\(^{50}\)

UTF8
  Eight-Bit Unicode Transformation Format\(^{51}\)

UUID
  Universally Unique Identifier

WAL
  Write-Ahead Log

XID
  Transaction Identifier

47. http://www.tpc.org/
50. http://www.unicode.org/
XML

Extensible Markup Language\textsuperscript{52}

\textsuperscript{52} http://en.wikipedia.org/wiki/XML
Bibliography

Selected references and readings for SQL and PostgreSQL.

Some white papers and technical reports from the original POSTGRES development team are available at the University of California, Berkeley, Computer Science Department web site¹.

SQL Reference Books


PostgreSQL-specific Documentation


Discusses SQL history and syntax, and describes the addition of INTERSECT and EXCEPT constructs into PostgreSQL. Prepared as a Master’s Thesis with the support of O. Univ. Prof. Dr. Georg Gottlob and Univ. Ass. Mag. Katrin Seyr at Vienna University of Technology.


Zelaine Fong, *The design and implementation of the POSTGRES query optimizer* ², University of California, Berkeley, Computer Science Department.

¹. http://db.cs.berkeley.edu/papers/
Proceedings and Articles


M. Stonebraker, L. A. Rowe, and M. Hirohama, “The implementation of POSTGRES”, *Transactions on Knowledge and Data Engineering 2*(1), IEEE, March 1990.


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Index

Symbols

$, 34
chtslibdir, 1004
chtslibdir/plugins, 563, 1672
*, 108
_pg_service.conf, 783
;:, 41
_PG_fini, 1004
_PG_init, 1004
_PG_output_plugin_init, 1282

A

abbrev, 253
ABORT, 1295
abs, 189
acos, 191
acosd, 191
administration tools
  externally maintained, 2524
adminpack, 2348
advisory lock, 426
age, 235
aggregate function, 11
  built-in, 290
  invocation, 37
  moving aggregate, 1028
  ordered set, 1031
  partial aggregation, 1032
  polymorphic, 1029
  support functions for, 1033
  user-defined, 1026
  variadic, 1029
AIX
  installation on, 470
  IPC configuration, 495
akeys, 2410
alias
  for table name in query, 10
  in the FROM clause, 98
  in the select list, 108
ALL, 301, 304
allow_system_table_mods configuration parameter, 571
ALTER AGGREGATE, 1297
ALTER COLLATION, 1299
ALTER CONVERSION, 1301
ALTER DATABASE, 1303
ALTER DEFAULT PRIVILEGES, 1306
ALTER DOMAIN, 1309
ALTER EVENT TRIGGER, 1313
ALTER EXTENSION, 1314
ALTER FOREIGN DATA WRAPPER, 1318
ALTER FOREIGN TABLE, 1320
ALTER FUNCTION, 1325
ALTER GROUP, 1329
ALTER INDEX, 1331
ALTER LANGUAGE, 1334
ALTER LARGE OBJECT, 1335
ALTER MATERIALIZED VIEW, 1336
ALTER OPERATOR, 1338
ALTER OPERATOR CLASS, 1340
ALTER OPERATOR FAMILY, 1342
ALTER POLICY, 1346
ALTER ROLE, 594, 1348
ALTER RULE, 1352
ALTER SCHEMA, 1353
ALTER SEQUENCE, 1354
ALTER SERVER, 1357
ALTER SYSTEM, 1359
ALTER TABLE, 1361
ALTER TABLESPACE, 1374
ALTER TEXT SEARCH CONFIGURATION, 1376
ALTER TEXT SEARCH DICTIONARY, 1378
ALTER TEXT SEARCH PARSER, 1380
ALTER TEXT SEARCH TEMPLATE, 1381
ALTER TRIGGER, 1382
ALTER TYPE, 1384
ALTER USER, 1388
ALTER USER MAPPING, 1389
ALTER VIEW, 1391
ANALYZE, 619, 1393
AND (operator), 185
anonymous code blocks, 1572
any, 183, 293, 301, 304
anyarray, 183
anypelement, 183
anynonarray, 183
anynonarray, 183
applicable role, 921
apply_name configuration parameter, 549
arbitrary precision numbers, 120
archive_cleanup_command recovery parameter, 666
archive_command configuration parameter, 534
archive_mode configuration parameter, 534
archive_timeout configuration parameter, 534
area, 250
armor, 2442
ARRAY, 42, 160
accessing, 162
constant, 161
constructor, 42
declaration, 160
determination of result type, 358
I/O, 168
modifying, 164
of user-defined type, 1036
searching, 167
array_agg, 290, 2415
array_append, 285
array_cat, 285
array_dims, 285
array_fill, 285
array_length, 285
array_lower, 285
array_ndims, 285
array_nulls configuration parameter, 566
array_position, 285
array_positions, 285
array_prepend, 285
array_remove, 285
array_replace, 285
array_to_json, 273
array_to_string, 285
array_to_tsvector, 256
array_upper, 285
ascii, 194
asin, 191
asind, 191
ASSERT
in PL/pgSQL, 1155
assertions
in PL/pgSQL, 1155
asynchronous commit, 712
AT TIME ZONE, 244
atan, 191
atan2, 191
atan2d, 191
atand, 191
authentication_timeout configuration parameter, 520
auth_delay, 2349
auth_delay.milliseconds configuration parameter, 2349
auto-increment
(see serial)
autocommit
bulk-loading data, 445
psql, 1877
autovacuum
configuration parameters, 555
general information, 623
autovacuum configuration parameter, 555
autovacuum_analyze_scale_factor configuration parameter, 556
autovacuum_analyze_threshold configuration parameter, 556
autovacuum_freeze_max_age configuration parameter, 556
autovacuum_max_workers configuration parameter, 556
autovacuum_multixact_freeze_max_age configuration parameter, 557
autovacuum_naptime configuration parameter, 557
autovacuum_vacuum_cost_delay configuration parameter, 557
autovacuum_vacuum_cost_limit configuration parameter, 557
autovacuum_vacuum_scale_factor configuration parameter, 556
autovacuum_vacuum_threshold configuration parameter, 556
autovacuum_work_mem configuration parameter, 524
auto_explain, 2349
auto_explain.log_analyze configuration parameter, 2350
auto_explain.log_buffers configuration parameter, 2350
auto_explain.log_format configuration parameter, 2350
auto_explain.log_min_duration configuration parameter, 2350
auto_explain.log_nested_statements configuration parameter, 2351
auto_explain.log_timing configuration parameter, 2350
auto_explain.log_triggers configuration parameter, 2350
auto_explain.log_verbose configuration parameter, 2350
auto_explain.sample_rate configuration parameter, 2351
avals, 2410
average, 290
avg, 290

B

B-tree
(see index)
backend_flush_after configuration parameter, 528
Background workers, 1278
backslash escapes, 26
backslash_quote configuration parameter, 566
backup, 328, 627
base type, 984
BASE_BACKUP, 2062
BEGIN, 1396
BETWEEN, 186
BETWEEN SYMMETRIC, 187
BGWORKER_BACKEND_DATABASE_CONNECTION, 1278
BGWORKER_SHMEM_ACCESS, 1278
bgwriter_delay configuration parameter, 526
bgwriter_flush_after configuration parameter, 527
bwiter_lru_maxpages configuration parameter, 527
bgwriter_lru_multiplier configuration parameter, 527
bigint, 30, 120
bignumeric, 123
binary data, 127
functions, 209
binary string
concatenation, 209
length, 210
bison, 457
bit string
constant, 29
data type, 147
bit strings
functions, 211
bitmap scan, 365, 539
bit_and, 290
bit_length, 192

bit_or, 290
BLOB
(see large object)
block_size configuration parameter, 569
bloom, 2352
bonjour configuration parameter, 519
bonjour_name configuration parameter, 519
Boolean
data type, 139
operators
(see operators, logical)
bool_and, 290
bool_or, 290
booting
starting the server during, 490
box, 251
box (data type), 144
BRIN
(see index)
brin_metapage_info, 2433
brin_page_items, 2434
brin_page_type, 2433
brin_remap_data, 2434
brin_summarize_new_values, 339
broadcast, 253
BSD Authentication, 591
btree_gin, 2355
btree_gist, 2355
btrim, 194, 210
bt_metap, 2432
bt_page_items, 2433
bt_page_stats, 2433
bytea, 127
bytea_output configuration parameter, 561

C

C, 728, 813
C++, 1026
canceling
SQL command, 763
cardinality, 285
CASCADE
with DROP, 87
foreign key action, 58
Cascading Replication, 643
CASE, 282
determination of result type, 358
case sensitivity
of SQL commands, 25
cast
  I/O conversion, 1430
cbrt, 189
ceil, 189
ceiling, 189
center, 250
Certificate, 590
char, 125
character, 125
character set, 562, 570, 610
character string
  concatenation, 192
  constant, 26
  data types, 125
  length, 192
character varying, 125
char_length, 192
check constraint, 52
CHECK OPTION, 1557
checkpoint, 714, 1398
checkpoint_completion_target configuration parameter, 533
checkpoint_flush_after configuration parameter, 533
checkpoint_timeout configuration parameter, 533
checkpoint_warning configuration parameter, 533
check_function_bodies configuration parameter, 559
chkpass, 2357
chr, 194
cid, 181
cidr, 146
circle, 145, 251
citext, 2358
client authentication, 575
  timeout during, 520
client_encoding configuration parameter, 562
client_min_messages configuration parameter, 557
clock_timestamp, 235
CLOSE, 1399
CLUSTER, 1401
  of databases
    (see database cluster)
clusterdb, 1766
clustering, 643
cluster_name configuration parameter, 554
cmax, 60
cmin, 60
COALESCE, 284
COLLATE, 41
collation, 607
  in PL/pgSQL, 1123
  in SQL functions, 1000
collation for, 316
column, 5, 50
  adding, 62
  removing, 62
  renaming, 64
  system column, 60
column data type
  changing, 63
column reference, 34
col_description, 322
COMMENT, 1404
  about database objects, 322
  in SQL, 32
COMMIT, 1408
COMMIT PREPARED, 1409
commit_delay configuration parameter, 533
commit_siblings configuration parameter, 533
common table expression
  (see WITH)
comparison
  composite type, 304
  operators, 185
  row constructor, 304
  subquery result row, 301
compiling
  libpq applications, 790
  composite type, 169, 984
  comparison, 304
  constant, 170
  constructor, 44
  computed field, 174
  concat, 194
  concat_ws, 194
  concurrency, 415
conditional expression, 282
configuration
  of recovery
    of a standby server, 666
    of the server, 513
    of the server
      functions, 326
configure, 458
config_file configuration parameter, 517
correlation, 185
connectby, 2490, 2497
connection service file, 783
conninfo, 734
constant, 26
constraint, 52
    adding, 62
    check, 52
    exclusion, 60
    foreign key, 57
    name, 52
    NOT NULL, 54
    primary key, 56
    removing, 63
    unique, 55
constraint exclusion, 83, 543
constraint_exclusion configuration parameter, 543
CONTINUE
    in PL/pgSQL, 1139
continuous archiving, 627
    in standby, 655
control file, 1057
convert, 194
convert_from, 194
convert_to, 194
COPY, 7, 1411
    with libpq, 766
corr, 293
correlation, 293
cos, 191
cosd, 191
cot, 191
cotd, 191
count, 290
covariance
    population, 293
    sample, 293
covar_pop, 293
covar_samp, 293
cpu_index_tuple_cost configuration parameter, 541
cpu_operator_cost configuration parameter, 541
cpu_tuple_cost configuration parameter, 541
CREATE ACCESS METHOD, 1421
CREATE AGGREGATE, 1423
CREATE CAST, 1430
CREATE COLLATION, 1435
CREATE CONVERSION, 1437
CREATE DATABASE, 599, 1439
CREATE DOMAIN, 1442
CREATE EVENT TRIGGER, 1445
CREATE EXTENSION, 1447
CREATE FOREIGN DATA WRAPPER, 1450
CREATE FOREIGN TABLE, 1452
CREATE FUNCTION, 1456
CREATE GROUP, 1465
CREATE INDEX, 1466
CREATE LANGUAGE, 1473
CREATE MATERIALIZED VIEW, 1476
CREATE OPERATOR, 1478
CREATE OPERATOR CLASS, 1481
CREATE OPERATOR FAMILY, 1484
CREATE POLICY, 1486
CREATE ROLE, 593, 1491
CREATE RULE, 1496
CREATE SCHEMA, 1499
CREATE SEQUENCE, 1502
CREATE SERVER, 1506
CREATE TABLE, 5, 1508
CREATE TABLE AS, 1523
CREATE TABLESPACE, 603, 1526
CREATE TEXT SEARCH CONFIGURATION, 1528
CREATE TEXT SEARCH DICTIONARY, 1530
CREATE TEXT SEARCH PARSER, 1532
CREATE TEXT SEARCH TEMPLATE, 1534
CREATE TRANSFORM, 1536
CREATE TRIGGER, 1539
CREATE USER, 1554
CREATE USER MAPPING, 1555
CREATE VIEW, 1557
createlang, 1773
createtable, 1776
CREATE_REPLICATION_SLOT, 2059
cross compilation, 464
cross join, 94
crosstab, 2491, 2493, 2494
crypt, 2439
cstring, 183
citid, 60, 1092
CUBE, 105
cube (extension), 2360
cume_dist, 299
    hypothetical, 297
current_catalog, 311
current_database, 311
current_date, 235
current_query, 311
current_role, 311
2551
current_schema, 311
current_schemas, 311
current_setting, 326
current_time, 235
current_timestamp, 235
current_user, 311
currval, 280
cursor
  CLOSE, 1399
  DECLARE, 1563
  FETCH, 1650
  in PL/pgSQL, 1147
  MOVE, 1676
  showing the query plan, 1645
cursor_tuple_fraction configuration parameter, 543
custom scan provider
  handler for, 2124
Cygwin
  installation on, 473

data area
  (see database cluster)
data partitioning, 643
data type, 118
  base, 984
  category, 349
  composite, 984
  constant, 30
  conversion, 348
  enumerated (enum), 141
  internal organization, 1005
  numeric, 119
  type cast, 41
  user-defined, 1033
database, 599
  creating, 2
  privilege to create, 594
database activity
  monitoring, 670
database cluster, 5, 488
data_checksums configuration parameter, 569
data_directory configuration parameter, 517
data_sync_retry configuration parameter, 568
date, 129, 131
  constants, 134
current, 245
  output format, 134
  (see also formatting)
  DateStyle configuration parameter, 561
date_part, 235, 240
date_trunc, 235, 244
dblink, 2366, 2372
dblink_build_sql_delete, 2396
dblink_build_sql_insert, 2394
dblink_build_sql_update, 2398
dblink_cancel_query, 2391
dblink_close, 2381
dblink_connect, 2367
dblink_connect_u, 2370
dblink_disconnect, 2371
dblink_error_message, 2384
dblink_exec, 2375
dblink_fetch, 2379
dblink_get_connections, 2383
dblink_get_notify, 2387
dblink_get_pkey, 2392
dblink_get_result, 2388
dblink_is_busy, 2386
dblink_open, 2377
dblink_send_query, 2385
db_user_namespace configuration parameter, 522
deadlock, 425
  timeout during, 565
deadlock_timeout configuration parameter, 565
DEALLOCATE, 1562
dearmor, 2442
debug_assertions configuration parameter, 569
debug_deadlocks configuration parameter, 572
debug_pretty_print configuration parameter, 549
debug_print_parse configuration parameter, 549
debug_print_plan configuration parameter, 549
debug_print_rewritten configuration parameter, 549
decimal
  (see numeric)
DECLARE, 1563
decode, 194, 210
decode_bytea
  in PL/Perl, 1195
decrypt, 2446
decrypt_iv, 2446

default value, 51
changing, 63

default_statistics_target configuration parameter, 543

default_tablespace configuration parameter, 558

default_text_search_config configuration parameter, 563
default_transaction_deferrable configuration parameter, 559
default_transaction_isolation configuration parameter, 559
default_transaction_read_only configuration parameter, 559
default_with_oids configuration parameter, 566
deferrable transaction setting, 1743
setting default, 559
defined, 2410
degrees, 189
delay, 247
DELETE, 13, 91, 1567, 2410
RETURNING, 91
deleting, 91
dense_rank, 299
hypothetical, 297
diameter, 250
dict_int, 2400
dict_xsyn, 2400
difference, 2406
digest, 2438
dirty read, 415
DISCARD, 1570
disjunction, 185
disk drive, 717
disk space, 618
disk usage, 708
DISTINCT, 8, 109
div, 189
dmetaphone, 2408
dmetaphone_alt, 2408
DO, 1572
document
text search, 376
dollar quoting, 28
double precision, 122
DROP ACCESS METHOD, 1574
DROP AGGREGATE, 1575
DROP CAST, 1577
DROP COLLATION, 1579

DROP CONVERSION, 1581
DROP DATABASE, 602, 1582
DROP DOMAIN, 1583
DROP EVENT TRIGGER, 1584
DROP EXTENSION, 1585
DROP FOREIGN DATA WRAPPER, 1587
DROP FOREIGN TABLE, 1589
DROP FUNCTION, 1591
DROP GROUP, 1593
DROP INDEX, 1594
DROP LANGUAGE, 1596
DROP MATERIALIZED VIEW, 1598
DROP OPERATOR, 1600
DROP OPERATOR CLASS, 1602
DROP OPERATOR FAMILY, 1604
DROP OWNED, 1606
DROP POLICY, 1608
DROP ROLE, 593, 1610
DROP RULE, 1612
DROP SCHEMA, 1614
DROP SEQUENCE, 1616
DROP SERVER, 1618
DROP TABLE, 6, 1620
DROP TABLESPACE, 1622
DROP TEXT SEARCH CONFIGURATION, 1624
DROP TEXT SEARCH DICTIONARY, 1626
DROP TEXT SEARCH PARSER, 1628
DROP TEXT SEARCH TEMPLATE, 1630
DROP TRANSFORM, 1632
DROP TRIGGER, 1634
DROP TYPE, 1636
DROP USER, 1638
DROP USER MAPPING, 1639
DROP VIEW, 1641
dropdb, 602, 1781
droplang, 1784
dropuser, 593, 1787
DROP_REPLICATION_SLOT, 2062
DTD, 152
DTrace, 465, 696
duplicate, 8
duplicates, 109
dynamic loading, 564, 1004
dynamic_library_path, 1004
dynamic_library_path configuration parameter, 564
dynamic_shared_memory_type configuration parameter, 524
each, 2410
earth, 2402
earthdistance, 2402
earth_box, 2402
earth_distance, 2402
ECPG, 813, 1790
effective_cache_size configuration parameter, 541
effective_io_concurrency configuration parameter, 527
elog, 2084
  in PL/Perl, 1195
  in PL/Python, 1217
  in PL/Tcl, 1183
embedded SQL
  in C, 813
enabled role, 945
enable_bitmapscan configuration parameter, 539
enable_hashagg configuration parameter, 539
enable_hashjoin configuration parameter, 539
enable_indexonlyscan configuration parameter, 539
enable_indexscan configuration parameter, 539
enable_material configuration parameter, 539
enable_mergejoin configuration parameter, 539
enable_nestloop configuration parameter, 540
enable_seqscan configuration parameter, 540
enable_sort configuration parameter, 540
enable_tidscan configuration parameter, 540
encode, 194, 210
  encode_array_constructor
    in PL/Perl, 1196
  encode_array_literal
    in PL/Perl, 1196
encode_bytea
  in PL/Perl, 1196
encode_typed_literal
  in PL/Perl, 1196
encrypt, 2446
encryption, 506
  for specific columns, 2437
encrypt_iv, 2446
END, 1642
enumerated types, 141
enum_first, 248
enum_last, 248
enum_range, 248
environment variable, 781
erreport, 2084
error codes
  libpq, 750
  list of, 2202
error message, 742
escape string syntax, 26
escape_string_warning configuration parameter, 567
escaping strings
  in libpq, 756
event log
  event log, 512
event trigger, 1075
    in C, 1081
    in PL/Tcl, 1185
event_source configuration parameter, 547
event_trigger, 183
every, 290
EXCEPT, 109
exceptions
  in PL/pgSQL, 1143
  in PL/Tcl, 1186
exclusion constraint, 60
EXECUTE, 1643
exist, 2410
EXISTS, 301
EXIT
  in PL/pgSQL, 1139
exit_on_error configuration parameter, 568
exp, 189
EXPLAIN, 431, 1645
expression
  order of evaluation, 45
  syntax, 33
extending SQL, 984
extension, 1055
  externally maintained, 2525
  external_pid_file configuration parameter, 517
extract, 235, 240
extra_float_digits configuration parameter, 562
gin_fuzzy_search_limit configuration parameter, 565
gin_leafpage_items, 2435
gin_metapage_info, 2434
gin_pageOpaque_info, 2434
gin_pending_list_limit configuration parameter, 561
GiST
(see index)
global data
  in PL/Python, 1210
  in PL/Tcl, 1181
GRANT, 64, 1654
GREATEST, 284
determination of result type, 358
Gregorian calendar, 2214
GROUP BY, 12, 102
grouping, 102, 298
GROUPING SETS, 105
GSSAPI, 584
GUID, 150

hash
(see index)
has_any_column_privilege, 314
has_column_privilege, 314
has_database_privilege, 314
has_foreign_data_wrapper_privilege, 314
has_function_privilege, 314
has_language_privilege, 314
has_schema_privilege, 314
has_sequence_privilege, 314
has_server_privilege, 314
has_tablespace_privilege, 314
has_table_privilege, 314
has_type_privilege, 314
HAVING, 12, 104
hba_file configuration parameter, 517
heap_page_items, 2432
heap_page_item attrs, 2432
height, 250
hierarchical database, 5
high availability, 643
history
  of PostgreSQL, xlv
 hmac, 2438
host, 253
host name, 736
hostmask, 253
Hot Standby, 643
hot_standby configuration parameter, 537
hot_standby_feedback configuration parameter, 538
HP-UX
  installation on, 474
  IPC configuration, 496
  shared library, 1014
hstore, 2408, 2410
hstore_to_array, 2410
hstore_to_json, 2410
hstore_to_jsonb, 2410
hstore_to_jsonb_loose, 2410
hstore_to_json_loose, 2410
hstore_to_matrix, 2410
huge_pages configuration parameter, 523
hypothetical-set aggregate
  built-in, 297
icount, 2417
ident, 586
identifier
  length, 25
  syntax of, 24
IDENTIFY_SYSTEM, 2059
ident_file configuration parameter, 517
idle_in_transaction_session_timeout configuration parameter, 560
idx, 2417
IFNULL, 284
ignore_checksum_failure configuration parameter, 573
ignore_system_indexes configuration parameter, 571
IMMUTABLE, 1001
IMPORT FOREIGN SCHEMA, 1661
IN, 301, 304
include
  in configuration file, 515
  in dir
    in configuration file, 516
include_if_exists
  in configuration file, 516
index, 360, 2427
  and ORDER BY, 364
B-tree, 361
BRIN, 363, 2176
building concurrently, 1469
combining multiple indexes, 365
examining usage, 373
on expressions, 366
for user-defined data type, 1043
GIN, 362, 2169
text search, 410
GiST, 361, 2149
text search, 410
hash, 361
index-only scans, 371
locks, 429
multicolumn, 363
partial, 367
SP-GiST, 362, 2160
unique, 366
index scan, 539
index-only scan, 371
index_am_handler, 183
inet (data type), 145
inet_client_addr, 312
inet_client_port, 312
inet_merge, 253
inet_same_family, 253
inet_server_addr, 312
inet_server_port, 312
information schema, 920
inheritance, 20, 75, 567
initcap, 194
initdb, 488, 1897
Initialization Fork, 2187
input function, 1034
INSERT, 6, 89, 1663
RETURNING, 91
inserting, 89
installation, 456
on Windows, 481
instr function, 1176
int, 31
(see smallint)
int4
(see integer)
int8
(see bigint)
intagg, 2415
intarray, 2417
integer, 30, 120
integer_datetimes configuration parameter, 569
interfaces
externally maintained, 2524
internal, 183
INTERSECT, 109
interval, 129, 137
output format, 139
(see also formatting)
IntervalStyle configuration parameter, 562
intset, 2417
int_array_aggregate, 2415
int_array_enum, 2415
inverse distribution, 295
IS DISTINCT FROM, 187, 304
IS DOCUMENT, 265
IS FALSE, 187
IS NOT DISTINCT FROM, 187, 304
IS NOT DOCUMENT, 265
IS NOT FALSE, 187
IS NOT NULL, 187
IS NOT TRUE, 187
IS NOT UNKNOWN, 187
IS NULL, 187, 568
IS TRUE, 187
IS UNKNOWN, 187
isclosed, 250
isempty, 290
isfinite, 235
isn, 2419
ISNULL, 187
isn_weak, 2421
isopen, 250
is_array_ref
in PL/Perl, 1196
is_valid, 2421
JOIN, 9, 94
controlling the order, 443
cross, 94
left, 95
natural, 95
outer, 10, 95
right, 95
self, 10
join_collapse_limit configuration parameter, 543
JSON, 153
functions and operators, 271
JSONB, 153
containment, 156
existence, 156
indexes on, 158
J
K

key word
  list of, 2216
  syntax of, 24

krb_caseins_users configuration parameter, 522

krb_server_keyfile configuration parameter, 521

L

label
  (see alias)
lag, 299
language_handler, 183
large object, 801
lastval, 280
last_value, 299

LATERAL
  in the FROM clause, 100
latitude, 2402
lca, 2427

lc_collate configuration parameter, 569
lc_ctype configuration parameter, 569
lc_messages configuration parameter, 562
lc_monetary configuration parameter, 562
lc_numeric configuration parameter, 563
lc_time configuration parameter, 563
LDAP, 461, 587
LDAP connection parameter lookup, 784
ldconfig, 469
lead, 299

LEAST, 284
  determination of result type, 358
left, 194
left join, 95
length, 194, 210, 250, 256
  of a binary string
    (see binary strings, length)
  of a character string
    (see character string, length)

length(tsvector), 388
levenshtein, 2407
levenshtein_less_equal, 2407
lex, 457
libedit, 456
libperl, 457
libpq, 728
  single-row mode, 762
libpq-fe.h, 728, 739
libpq-int.h, 739
libpython, 457
library finalization function, 1004
library initialization function, 1004
LIKE, 212
   and locales, 606
LIMIT, 111
line, 143
line segment, 144
linear regression, 293
Linux
   IPC configuration, 496
   shared library, 1014
   start script, 491
LISTEN, 1670
listen_addresses configuration parameter, 518
ll_to_earth, 2402
ln, 189
lo, 2423
LOAD, 1672
load balancing, 643
locale, 489, 605
localtime, 235
localtimestamp, 235
local_preload_libraries configuration parameter, 563
lock, 421, 421, 1673
   advisory, 426
   monitoring, 694
lock_timeout configuration parameter, 560
log, 189
log shipping, 643
logging_collector configuration parameter, 545
Logical Decoding, 1282, 1284
login privilege, 594
log Autovacuum_min_duration configuration parameter, 555
log btree_build_stats configuration parameter, 572
log checkpoints configuration parameter, 549
log_connections configuration parameter, 549
log_destination configuration parameter, 544
log_directory configuration parameter, 545
log_disconnections configuration parameter, 550
log_duration configuration parameter, 550
log_errorverbosity configuration parameter, 550
log_executor_stats configuration parameter, 555
log_filename configuration parameter, 545
log_file_mode configuration parameter, 546
log_hostname configuration parameter, 550
log_line_prefix configuration parameter, 550
log_lock_waits configuration parameter, 551
log_min_duration_statement configuration parameter, 548
log_min_error_statement configuration parameter, 547
log_min_messages configuration parameter, 547
log_parser_stats configuration parameter, 555
log_planner_stats configuration parameter, 555
log_replication_commands configuration parameter, 552
log_rotation_age configuration parameter, 546
log_rotation_size configuration parameter, 546
log_statement configuration parameter, 552
log_statement_stats configuration parameter, 555
log_replication_commands configuration parameter, 552
log_temp_files configuration parameter, 552
log_timezone configuration parameter, 552
log_truncate_on_rotation configuration parameter, 546
long, 2402
looks_like_number
   in PL/Perl, 1196
loop
   in PL/pgSQL, 1138
lower, 192, 290
   and locales, 606
lower_inc, 290
lower_inf, 290
lo_close, 805
lo_compatible_privileges configuration parameter, 567
lo_creat, 802, 806
lo_create, 802
lo_export, 803, 806
lo_from_bytea, 806
lo_get, 806
lo_import, 802, 806
lo_import_with_oid, 802
lo_lseek, 804
MAC address
(see macaddr)
macaddr (data type), 147
macOS
installation on, 475
magic block, 1004
maintenance, 617
maintenance_work_mem configuration parameter, 524
make, 456
make_date, 235
make_interval, 235
make_time, 235
make_timestamp, 235
make_timestamptz, 235
make_valid, 2421
MANPATH, 470
masklen, 253
materialized view
implementation through rules, 1094
materialized views, 2028
max, 290
max_connections configuration parameter, 518
max_files_per_process configuration parameter, 525
max_function_args configuration parameter, 569
max_identifier_length configuration parameter, 570
max_index_keys configuration parameter, 570
max_locks_per_transaction configuration parameter, 565
max_parallel_workers_per_gather configuration parameter, 528
max_pred_locks_per_transaction configuration parameter, 566
max_prepared_transactions configuration parameter, 523
max_replication_slots configuration parameter, 535
max_stack_depth configuration parameter, 524
max_standby_archive_delay configuration parameter, 537
max_standby_streaming_delay configuration parameter, 538
max_wal_senders configuration parameter, 535
max_wal_size configuration parameter, 534
max_worker_processes configuration parameter, 534
md5, 194, 210, 583
median, 38
(see also percentile)
memory context
in SPI, 1263
memory overcommit, 500
metaphone, 2407
min, 290
MinGW
installation on, 476
min_parallel_relation_size configuration parameter, 541
min_wal_size configuration parameter, 534
mod, 189
mode
statistical, 295
monitoring
database activity, 670
MOVE, 1676
moving-aggregate mode, 1028
Multiversion Concurrency Control, 415
MultiXactId, 623
MVCC, 415
name
  qualified, 71
  syntax of, 24
  unqualified, 72
NaN
  (see not a number)
natural join, 95
negation, 185
NetBSD
  IPC configuration, 496
  shared library, 1014
  start script, 491
netmask, 253
network, 253
  data types, 145
Network Attached Storage (NAS)
  (see Network File Systems)
Network File Systems, 489
nextval, 280
NFS
  (see Network File Systems)
nlevel, 2427
  non-durable, 448
  nonblocking connection, 730, 758
  nonrepeatable read, 415
  normal_rand, 2490
NOT (operator), 185
  not a number
    double precision, 122
    numeric (data type), 121
NOT IN, 301, 304
  not-null constraint, 54
notation
  functions, 47
  notice processing
    in libpq, 774
notice processor, 774
notice receiver, 774
NOTIFY, 1678
  in libpq, 765
NOTNULL, 187
now, 235
npoints, 250
nth_value, 299
ntile, 299
null value
  with check constraints, 54
  comparing, 187
  default value, 51
  in DISTINCT, 109
  in libpq, 754
  in PL/Perl, 1189
  in PL/Python, 1206
    with unique constraints, 56
NULLIF, 284
number
  constant, 29
  numeric, 30
  numeric (data type), 120
numnode, 256, 389
num_nonnulls, 188
num_nulls, 188
NVL, 284
object identifier
  data type, 181
object-oriented database, 5
obj_description, 322
octet_length, 192, 209
OFFSET, 111
oid, 181
  column, 60
    in libpq, 755
oid2name, 2513
old_snapshot_threshold configuration parameter, 528
ON CONFLICT, 1663
ONLY, 94
OOM, 500
opaque, 183
OpenBSD
  IPC configuration, 496
  shared library, 1014
  start script, 491
OpenSSL, 461
  (see also SSL)
operator, 185
  invocation, 36
  logical, 185
  precedence, 32
  syntax, 31
    type resolution in an invocation, 349
    user-defined, 1038
operator class, 369, 1043
operator family, 369, 1051
operator_precedence_warning configuration parameter, 567
OR (operator), 185
Oracle
- porting from PL/SQL to PL/pgSQL, 1169
ORDER BY, 8, 110
- and locales, 606
ordered-set aggregate, 37
- built-in, 295
ordering operator, 1053
ordinality, 309
OS X
- IPC configuration, 497
- shared library, 1014
outer join, 95
output function, 1034
OVER clause, 39
overcommit, 500
OVERLAPS, 239
overlay, 192, 209
overloading
- functions, 1000
- operators, 1038
owner, 64

pageinspect, 2431
page_header, 2432
palloc, 1013
PAM, 461, 590
parallel query, 449
parallel_setup_cost configuration parameter, 541
parallel_tuple_cost configuration parameter, 541
parameter
- syntax, 34
parenthesis, 34
parse_ident, 194
partitioning, 78
password, 594
- authentication, 583
- of the superuser, 489
password file, 783
passwordcheck, 2435
password_encryption configuration parameter, 521
path, 251, 469
- for schemas, 557
path (data type), 144
pattern matching, 212
patterns
- in psql and pg_dump, 1876
pclose, 250
peer, 587
percentile
- continuous, 295
- discrete, 295
percent_rank, 299
- hypothetical, 297
performance, 431
perl, 458, 1188
permission
- (see privilege)
pfree, 1013
PGAPPNAME, 782
pgbench, 1800
PGcancel, 763
PGCLIENTENCODING, 782
PGconn, 728
PGCONNECT_TIMEOUT, 782
pgcrypto, 2437
PGDATA, 488
PGDATABASE, 782
PGDATESTYLE, 783
PGEventProc, 777
PGGEOQ, 783
PGGSLLIB, 782
PGHOST, 782
PGHOSTADDR, 782
PGKRBSRVNAME, 782
PGLOCALEDIR, 783
PGOPTIONS, 782
PGPASSFILE, 782
PGPASSWORD, 782
PGPORT, 782
pgp_armor_headers, 2443
pgp_key_id, 2442
pgp_pub_decrypt, 2442
pgp_pub_decrypt_bytea, 2442
pgp_pub_encrypt, 2442
pgp_pub_encrypt_bytea, 2442
pgp_sym_decrypt, 2441
pgp_sym_decrypt_bytea, 2441
pgp_sym_encrypt, 2441
pgp_sym_encrypt_bytea, 2441
PGREQUIREPEER, 782
PGREQUIRESSL, 782
PGresolt, 748
pgrowlocks, 2451, 2452
PGSERVICE, 782
PGSERVICEFILE, 782
PGSSLCERT, 782
PGSSLCOMPRESSION, 782
pg_get_object_address, 321
pg_get_ruledef, 316
pg_get_serial_sequence, 316
pg_get_triggerdef, 316
pg_get_userbyid, 316
pg_get_viewdef, 316
pg_group, 2024
pg_has_role, 314
pg_hba.conf, 575
pg_ident.conf, 582
pg_identify_object, 321
pg_identify_object_as_address, 321
pg_index, 1981
pg_index_has_property, 316
pg_indexes, 2024
pg_indexes_size, 337
pg_index_column_has_property, 316
pg_language, 1985
pg_largeobject, 1987
pg_largeobject_metadata, 1987
pg_last_committed_xact, 324
pg_last_xact_replay_timestamp, 330
pg_locks, 2025
pg_logdir_ls, 2348
pg_logical_emit_message, 333
pg_logical_slot_get_binary_changes, 333
pg_logical_slot_get_changes, 333
pg_logical_slot_peek_binary_changes, 333
pg_logical_slot_peek_changes, 333
pg_lsn, 183
pg_ls_dir, 341
pg_matviews, 2028
pg_my_temp_schema, 312
pg_namespace, 1988
pg_notification_queue_usage, 312
pg_notify, 1679
pg_opclass, 1988
pg_opclass_is_visible, 316
pg_opfamily, 1990
pg_opfamily_is_visible, 316
pg_options_to_table, 316
pg_pltemplate, 1991
pg_policies, 2029
pg_policy, 1991
pg_postmaster_start_time, 312
pg_prepared_statements, 2030
pg_prepared_xacts, 2031
pg_prewarm, 2451
pg_proc, 1992
pg_range, 1997
pg_read_binary_file, 341
pg_read_file, 341
pg_receivevexlog, 1837
pg_recvlogical, 1841
pg_relation_filenode, 339
pg_relation_filepath, 339
pg_relation_size, 337
pg_reload_conf, 327
pg_relpages, 2460
pg_replication_origin, 1998
pg_replication_origin_advance, 333
pg_replication_origin_create, 333
pg_replication_origin_drop, 333
pg_replication_origin_oid, 333
pg_replication_origin_progress, 333
pg_replication_origin_session_is_setup, 333
pg_replication_origin_session_progress, 333
pg_replication_origin_session_reset, 333
pg_replication_origin_xact_reset, 333
pg_replication_origin_xact_setup, 333
pg_replication_origin_xact_setup, 333
pg_replication_slots, 2032
pg_resetxlog, 1910
pg_restore, 1845
pg_rewind, 1913
pg_rewrite, 1999
pg_roles, 2034
pg_rotate_logfile, 327
pg_rules, 2035
pg_seclabel, 1999
pg_seclabels, 2035
pg_service.conf, 783
pg_settings, 2036
pg_shadow, 2038
pg_shdepend, 2000
pg_shdescription, 2001
pg_shseclabel, 2002
pg_size_bytes, 337
pg_size_pretty, 337
pg_sleep, 247
pg_sleep_for, 247
pg_sleep_until, 247
pg_standby, 2520
pg_start_backup, 328
pg_statio_all_indexes, 673
pg_statio_all_sequences, 673
pg_statio_all_tables, 673
pg_statio_sys_indexes, 673
pg_statio_sys_sequences, 673
pg_statio_sys_tables, 673
pg_statio_user_indexes, 673
pg_statio_user_sequences, 673
pg_statio_user_tables, 673
pg_statistic, 442, 2003
pg_stats, 442, 2039
pg_stat_activity, 672
pg_stat_all_indexes, 673
pg_stat_all_tables, 673
pg_stat_archiver, 673
pg_stat_bgwriter, 673
pg_stat_clear_snapshot, 692
pg_stat_database, 673
pg_stat_database_conflicts, 673
pg_stat_file, 341
pg_stat_get_activity, 692
pg_stat_get_snapshot_timestamp, 692
pg_stat_progress_vacuum, 672
pg_stat_replication, 672
pg_stat_reset, 692
pg_stat_reset_shared, 692
pg_stat_reset_single_function_counters, 692
pg_stat_reset_single_table_counters, 692
pg_stat_ssl, 672
pg_stat_statements, 2453
  function, 2456
pg_stat_statements_reset, 2456
pg_stat_sys_indexes, 673
pg_stat_sys_tables, 673
pg_stat_user_functions, 673
pg_stat_user_indexes, 673
pg_stat_user_tables, 673
pg_stop_backup, 328
pg_switch_xlog, 328
pg_tables, 2042
pg_tablespace, 2005
pg_tablespace_databases, 316
pg_tablespace_location, 316
pg_tablespace_size, 337
pg_table_is_visible, 316
pg_table_size, 337
pg_temp, 558
  securing functions, 1462
pg_terminate_backend, 327
pg_test_fsync, 1916
pg_test_timing, 1918
pg_timezone_abbrevs, 2043
pg_timezone_names, 2043
pg_total_relation_size, 337
pg_transform, 2005
pg_trgm, 2461
  pg_trgm.similarity_threshold configuration parameter, 2464
  pg_trgm.word_similarity_threshold configuration parameter, 2464
pg_trigger, 2006
pg_try_advisory_lock, 343
pg_try_advisory_lock_shared, 343
pg_try_advisory_xact_lock, 343
pg_try_advisory_xact_lock_shared, 343
pg_ts_config, 2008
  pg_ts_config_is_visible, 316
  pg_ts_config_map, 2008
  pg_ts_dict, 2009
  pg_ts_dict_is_visible, 316
pg_ts_parser, 2009
  pg_ts_parser_is_visible, 316
pg_ts_template, 2010
  pg_ts_template_is_visible, 316
pg_type, 2011
  pg_typeof, 316
  pg_type_is_visible, 316
pg_upgrade, 1922
pg_user, 2044
  pg_user_mapping, 2044
pg_views, 2045
pg_xact_commit_timestamp, 324
pg_xact_commit_timestamp, 324
pg_xact_notify, 1929
pg_xact_notifier, 328
pg_xact_notifier_offset, 328
pg_xact_location, 328
pg_xact_notification, 328
pg_xact_notification_offset, 328
pg_xact_rollback, 328
pg_xact_rollback_resume, 331
pg_xact_rollback_resume, 331
phantom read, 415
phraseto_tsquery, 256, 383
Pi, 189
PIC, 1013
PID
determining PID of server process
in libpq, 742
PITR, 627
PITR standby, 643
pkg-config
with ecpg, 870
with libpq, 790
PL/Perl, 1188
PL/PerlU, 1198
PL/pgSQL, 1116
PL/Python, 1203
PL/SQL (Oracle)
porting to PL/pgSQL, 1169
point, 143, 251
point-in-time recovery, 627
policy, 65
polygon, 145, 251
polymorphic function, 985
polymorphic type, 985
popen, 250
populate_record, 2410
port, 736
port configuration parameter, 518
position, 192, 209
POSTGRES, xliv, 1, 490, 599, 1932
postgres user, 488
Postgres95, xliv
postgresql.auto.conf, 514
postgresql.conf, 513
postgres_fdw, 2468
postmaster, 1940
post_auth_delay configuration parameter, 571
power, 189
PQbackendPID, 742
PQbinaryTuples, 753
with COPY, 767
PQcancel, 764
PQclear, 751
PQclientEncoding, 770
PQcmdStatus, 755
PQcmdTuples, 755
PQconfdefaults, 732
PQconnectdb, 729
PQconnectdbParams, 729
PQconnectionNeedsPassword, 742
PQconnectionUsedPassword, 742
PQconnectPoll, 730
PQconnectStart, 730
PQconnectStartParams, 730
PQconninfo, 732
PQconninfoFree, 772
PQconninfoParse, 732
PQconsumeInput, 761
PQcopyResult, 773
PQdb, 740
PQdescribePortal, 747
PQdescribePrepared, 747
PQencryptPassword, 772
PQendcopy, 770
PQerrorMessage, 742
PQescapeBytea, 758
PQescapeByteaConn, 757
PQescapeIdentifier, 756
PQescapeLiteral, 756
PQescapeString, 757
PQescapeStringConn, 757
PQexec, 744
PQexecParams, 744
PQexecPrepared, 747
PQfformat, 753
PQfsize, 753
PQftable, 752
PQftablecol, 753
PQftype, 753
PQgetCancel, 763
PQgetCopyData, 768
PQgetisnull, 754
with COPY, 767
PQfinish, 733
PQfireResultCreateEvents, 772
PQflush, 762
PQfmod, 752
PQfnumber, 752
2566
externally maintained, 2525
handler for, 2100
protocol
frontend-backend, 2046
ps
to monitor activity, 670
psql, 3, 1854
Python, 1203

qualified name, 71
query, 7, 93
query plan, 431
query tree, 1085
querytree, 256, 389
quotation marks
and identifiers, 25
escaping, 26
quote_all_identifiers configuration parameter, 567
quote_ident, 194
in PL/Perl, 1195
use in PL/pgSQL, 1129
quote_literal, 194
in PL/Perl, 1195
use in PL/pgSQL, 1129
quote_nullable, 194
in PL/Perl, 1195
use in PL/pgSQL, 1129

radians, 189
radius, 250, 589
RAISE
in PL/pgSQL, 1153
random, 191
random_page_cost configuration parameter, 540
range table, 1085
range type, 175
exclude, 180
indexes on, 180
rank, 299
hypothetical, 297
read committed, 416
read-only transaction
setting, 1743
regular expressions
and locales, 606
reindex, 625, 1690
reindexdb, 1889
relation, 5
relational database, 5
RELEASE SAVEPOINT, 1693
repeat, 194
repeatable read, 418
replace, 194
replacement_sort_tuples configuration parameter, 524
replication, 643
Replication Origins, 1291
Replication Progress Tracking, 1291
replication slot
logical replication, 1284
streaming replication, 651
reporting errors
in PL/pgSQL, 1153
RESET, 1695
restartpoint, 715
restart_after_crash configuration parameter, 568
restore_command recovery parameter, 666
RESTRICT
with DROP, 87
foreign key action, 58
RETURN NEXT
in PL/pgSQL, 1134
RETURN QUERY
in PL/pgSQL, 1134
RETURNING, 91
RETURNING INTO
in PL/pgSQL, 1126
reverse, 194
REVOKE, 64, 1697
right, 194
right join, 95
role, 593, 597
applicable, 921
enabled, 945
membership in, 595
privilege to create, 594
privilege to initiate replication, 594
ROLLBACK, 1701
psql, 1879
ROLLBACK PREPARED, 1702
ROLLBACK TO SAVEPOINT, 1704
ROLLUP, 105
round, 189
row, 5, 44, 50
row estimation
planner, 2194
row type, 169
constructor, 44
row-level security, 65
row-wise comparison, 304
row_number, 299
row_security configuration parameter, 558
row_security_active, 314
row_to_json, 273
rpad, 194
rtrim, 194
rule, 1085
and materialized views, 1094
and views, 1087
for DELETE, 1096
for INSERT, 1096
for SELECT, 1087
compared with triggers, 1109
for UPDATE, 1096
SAVEPOINT, 1706
savepoints
defining, 1706
releasing, 1693
rolling back, 1704
scalar
(see expression)
scale, 189
schema, 70, 599
creating, 70
current, 72, 311
public, 71
removing, 71
SCO
installation on, 476
SCO OpenServer
IPC configuration, 497
search path, 72
current, 311
object visibility, 315
search_path configuration parameter, 72, 557
use in securing functions, 1462
SECURITY LABEL, 1708
sec_to_gc, 2402
SPI_copytuple, 1267
SPI_cursor_close, 1252
    in PL/Perl, 1193
SPI_cursor_fetch, 1248
SPI_cursor_find, 1247
SPI_cursor_move, 1249
SPI_cursor_open, 1242
SPI_cursor_open_with_args, 1244
SPI_cursor_open_with_paramlist, 1246
SPI_exec, 1228
SPI_execp, 1241
SPI_execute, 1225
SPI_execute_plan, 1238
SPI_execute_plan_with_paramlist, 1240
SPI_execute_with_args, 1229
spi_exec_prepared
    in PL/Perl, 1194
spi_exec_query
    in PL/Perl, 1192
spi_fetchrow
    in PL/Perl, 1193
SPI_finish, 1222
SPI_fname, 1255
SPI_fnumber, 1256
SPI_freeplan, 1273
    in PL/Perl, 1194
SPI_freetuple, 1271
SPI_freetuptable, 1272
SPI_getargcount, 1235
SPI_getargtypeid, 1236
SPI_getbinval, 1258
SPI_getnspname, 1262
SPI_getrelname, 1261
SPI_gettype, 1259
SPI_gettypeid, 1260
SPI_getvalue, 1257
SPI_is_cursor_plan, 1237
SPI_kEEPplan, 1253
spi_lastoid
    in PL/Tcl, 1183
SPI_modifytuple, 1269
SPI_pallocate, 1263
SPI_pfree, 1266
SPI_pop, 1224
SPI_prepare, 1231
    in PL/Perl, 1194
SPI_prepare_cursor, 1233
SPI_prepare_params, 1234
SPI_push, 1223
spi_query
    in PL/Perl, 1193
spi_query_prepared
    in PL/Perl, 1194
SPI_repalloc, 1265
SPI_returntuple, 1268
SPI_saveplan, 1254
SPI_scroll_cursor_fetch, 1250
SPI_scroll_cursor_move, 1251
split_part, 194
SQL/CLI, 2239
SQL/Foundation, 2239
SQL/Framework, 2239
SQL/JRT, 2239
SQL/MED, 2239
SQL/OLB, 2239
SQL/PSM, 2239
SQL/Schemata, 2239
SQL/XML, 2239
sql_inheritance configuration parameter, 567
sqrt, 189
ssh, 511
SSL, 415
SSL, 507, 785
    with libpq, 738, 744
ssl configuration parameter, 520
sslinfo, 2488
ssl_ca_file configuration parameter, 520
ssl_cert_file configuration parameter, 520
ssl_cipher, 2488
ssl_ciphers configuration parameter, 521
ssl_client_cert_present, 2488
ssl_client_dn, 2488
ssl_client_dn_field, 2489
ssl_client_serial, 2488
ssl_crl_file configuration parameter, 520
ssl_ecdh_curve configuration parameter, 521
ssl_extension_info, 2489
ssl_issuer_dn, 2488
ssl_issuer_field, 2489
ssl_is_used, 2488
ssl_key_file configuration parameter, 520
ssl_prefer_server_ciphers configuration parameter, 521
ssl_version, 2488
SSPI, 585
STABLE, 1001
standard deviation, 293
    population, 293
    sample, 293
standard_conforming_strings configuration parameter, 567
standby server, 643
standby_mode recovery parameter, 668
START TRANSACTION, 1748
START_REPLICATION, 2059
statement_timeout configuration parameter, 559
statement_timestamp, 235
statistics, 293, 671
of the planner, 441, 619
stats_temp_directory configuration parameter, 555
stddev, 293
stddev_pop, 293
stddev_samp, 293
STONITH, 643
storage parameters, 1515
Streaming Replication, 643
string
(see character string)
strings
backslash quotes, 566
escape warning, 567
standard conforming, 567
string_agg, 290
string_to_array, 285
strip, 256, 388
strpos, 194
substring, 192, 209, 213, 214
sum, 290
superuser, 4, 594
superuser_reserved_connections configuration parameter, 518
suppress_redundant_updates_trigger, 343
svals, 2410
synchronize_seqscans configuration parameter, 568
synchronous commit, 712
Synchronous Replication, 643
synchronous_commit configuration parameter, 530
synchronous_standby_names configuration parameter, 536
syntax
SQL, 24
syslog_facility configuration parameter, 546
syslog_ident configuration parameter, 547
syslog_sequence_numbers configuration parameter, 547
table, 5, 50
creating, 50
inheritance, 75
modifying, 61
partitioning, 78
removing, 51
renaming, 64
TABLE command, 1711
table expression, 93
table function, 99
table sampling method, 2121
tablefunc, 2489
tableoid, 60
TABLESAMPLE method, 2121
tablespace, 602
default, 558
temporary, 558
tan, 191
tand, 191
target list, 1086
Tcl, 1179
tcn, 2499
tcp_keepalives_count configuration parameter, 520
tcp_keepalives_idle configuration parameter, 519
tcp_keepalives_interval configuration parameter, 519
template0, 600
template1, 600, 600
temp_buffers configuration parameter, 523
temp_file_limit configuration parameter, 525
temp_tablespaces configuration parameter, 558
test, 718
test_decoding, 2500
text, 125, 253
text search, 375
data types, 148
functions and operators, 148
2572
indexes, 410
text2ltree, 2427
threads
with libpq, 789
tid, 181
time, 129, 132
constants, 134
current, 245
output format, 134
(see also formatting)
time span, 129
time with time zone, 129, 132
time without time zone, 129, 132
time zone, 135, 562
conversion, 244
input abbreviations, 2213
time zone data, 464
time zone names, 562
timelines, 627
TIMELINE_HISTORY, 2059
timeofday, 235
timeout
client authentication, 520
deadlock, 565
timestamp, 129, 133
timestamp with time zone, 129, 133
timestamp without time zone, 129, 133
timestamptz, 129
TimeZone configuration parameter, 562
timezone_abbreviations configuration parameter, 562
TOAST, 2183
and user-defined types, 1036
per-column storage settings, 1363
versus large objects, 801
token, 24
to_ascii, 194
to_char, 227
and locales, 607
to_date, 227
to_hex, 194
to_json, 273
to_jsonb, 273
to_timestamp, 227, 235
to_tsquery, 256, 382
to_tsvector, 256, 381
trace_locks configuration parameter, 571
trace_lock_oidmin configuration parameter, 572
trace_lock_table configuration parameter, 572
trace_lwlocks configuration parameter, 572
trace_notify configuration parameter, 571
trace_recovery_messages configuration parameter, 571
trace_sort configuration parameter, 571
trace_userlocks configuration parameter, 572
track_activities configuration parameter, 554
track_activity_query_size configuration parameter, 554
track_commit_timestamp configuration parameter, 536
track_counts configuration parameter, 554
track_functions configuration parameter, 555
track_io_timing configuration parameter, 554
transaction, 15
transaction ID
wraparound, 620
transaction isolation, 415
transaction isolation level, 416
read committed, 416
repeatable read, 418
serializable, 419
setting, 1743
setting default, 559
transaction log
(see WAL)
transaction_timestamp, 235
transform_null_equals configuration parameter, 568
translate, 194
trigger, 183, 1066
arguments for trigger functions, 1068
for updating a derived tsvector column, 391
in C, 1069
in PL/pgSQL, 1155
in PL/Python, 1211
in PL/Tcl, 1183
compared with rules, 1109
triggered_change_notification, 2499
trigger_file recovery parameter, 669
2573
trim, 192, 209
true, 139
trup, 189, 255
TRUNCATE, 1749
trusted
PL/Perl, 1197
tsearch2, 2500
tsm_handler, 183
tsm_system_rows, 2502
tsm_system_time, 2503
tquery (data type), 149
tquery_phrase, 256, 389
tsvector (data type), 148
tsvector_concatenation, 388
tsvector_to_array, 256
hsvector_update_trigger, 256
tvector_update_trigger_column, 256
ts_debug, 259, 406
t_delete, 256
ts_filter, 256
ts_headline, 256, 386
ts_lexize, 259, 409
ts_parse, 259, 408
ts_rank, 256, 384
ts_rank_cd, 256, 384
ts_rewrite, 256, 389
ts_stat, 259, 392
ts_token_type, 259, 408
tuple_data_split, 2432
tsid_current, 323
tsid_current_snapshot, 323
tsid_snapshot_xip, 323
tsid_snapshot xmax, 323
tsid_snapshot xmin, 323
tsid_visible_in_snapshot, 323
type
(see data type)
polymorphic, 985
type cast, 30, 41
unique constraint, 55
Unix domain socket, 736
UNIX
installation on, 476
IPC configuration, 498
shared library, 1015
unix_socket_directories configuration parameter, 518
unix_socket_group configuration parameter, 518
unix_socket_permissions configuration parameter, 519
UNLISTEN, 1752
unnest, 285
for tsvector, 256
unqualified name, 72
updatable views, 1559
UPDATE, 12, 90, 1754
RETURNING, 91
update_process_title configuration parameter, 554
upgrading, 90
upgrading, 503
upper, 192, 290
and locales, 606
upper inc, 290
upper inf, 290
UPSERT, 1663
URI, 734
user, 311, 593
current, 311
user mapping, 86
User name maps, 581
UUID, 150, 462
uuid-ossp, 2505
uuid_generate_v1, 2505
uuid_generate_v1mc, 2505
uuid_generate_v3, 2505
uuid_generate_v4, 2505
UESCAPE, 25, 28
unaccent, 2503, 2505
Unicode escape
in identifiers, 25
in string constants, 28
UNION, 109
determination of result type, 358
uniq, 2417
vacuum, 617, 1759
vacuumdb, 1892
vacuumlo, 2518
vacuum_cost_delay configuration parameter, 525
vacuum_cost_limit configuration parameter, 526
vacuum_cost_page_dirty configuration parameter, 526
vacuum_cost_page_hit configuration parameter, 526
vacuum_cost_page_miss configuration parameter, 526
vacuum_defer_cleanup_age configuration parameter, 537
vacuum_freeze_min_age configuration parameter, 560
vacuum_freeze_table_age configuration parameter, 560
vacuum_multixact_freeze_min_age configuration parameter, 561
vacuum_multixact_freeze_table_age configuration parameter, 560
value expression, 33
VALUES, 111, 1762
determination of result type, 358
varchar, 125
variadic function, 993
variance, 293
population, 293
sample, 293
var_pop, 293
var_samp, 293
version, 4, 312
compatibility, 503
view, 14
implementation through rules, 1087
materialized, 1094
updating, 1101
Visibility Map, 2186
VM
(see Visibility Map)
void, 183
VOLATILE, 1001
volatility
functions, 1001
VPATH, 458, 1064
wal_level configuration parameter, 529
wal_log_hints configuration parameter, 532
wal_receiver_status_interval configuration parameter, 538
wal_receiver_timeout configuration parameter, 538
wal_retrieve_retry_interval configuration parameter, 539
wal_segment_size configuration parameter, 570
wal_sender_timeout configuration parameter, 536
wal_sync_method configuration parameter, 531
wal_writer_delay configuration parameter, 532
wal_writer_flush_after configuration parameter, 532
warm standby, 643
WHERE, 101
where to log, 544
WHILE
in PL/pgSQL, 1140
width, 250
width_bucket, 189
window function, 17
built-in, 299
invocation, 39
order of execution, 107
WITH
in SELECT, 112, 1711
WITH CHECK OPTION, 1557
WITHIN GROUP, 37
witness server, 643
word_similarity, 2462
work_mem configuration parameter, 523
wraparound
of multixact IDs, 623
of transaction IDs, 620
xid, 181
xmax, 60
xmin, 60
XML, 151
XML export, 267
XML option, 152, 561
xml2, 2507
xmlagg, 264, 290
xmlbinary configuration parameter, 561
xmlcomment, 260
xmlconcat, 261
xmlelement, 261
XMLEXISTS, 265
xmlforest, 263
xmloption configuration parameter, 561
xmlparse, 151
xmlpi, 263
xmlroot, 263
xmlserialize, 152
xml_is_well_formed, 265
xml_is_well_formed_content, 265
xml_is_well_formed_document, 265
XPath, 266
xpath_exists, 267
xpath_table, 2509
xslt_process, 2511

Y

yacc, 457

Z

zero_damaged_pages configuration parameter, 573
zlib, 457, 464